

LL(1) Parsing

Motivation

- ▶ We usually need a parse tree
- ▶ Now we have all the machinery we need to produce it!

Parsing

- ▶ Create a stack
- ▶ This keeps track of the symbols that we have yet to match up with input tokens
- ▶ Initially, push start symbol to stack
- ▶ Note: We assume we have pseudotoken “\$” appended to the input token list

Parse

► Now we cycle through a loop:

```
1 while input remains:
2     let t = symbol of next token of input
3     let s = symbol at top of stack
4     if s is a terminal:
5         if s == t:
6             consume next token
7             pop stack
8     else:
9         pop stack
10        push symbols in table[s][t] in reverse order
```

Parse

- ▶ Problem: This algorithm has no error checking
- ▶ We probably want to check for a few error conditions...

Errors

- ▶ If we come into the while-loop and stack is empty: Error
 - ▶ Indicates we have trailing garbage at end of input

Errors

- ▶ If we exit the while-loop and stack is not empty: Error
 - ▶ Indicates we had early EOF

Errors

- ▶ If s is a terminal but $s \neq t$: Error
 - ▶ We expected to see symbol s , but we actually saw symbol t

Errors

- ▶ If s is a nonterminal and $\text{table}[s][t]$ is empty: Error
 - ▶ Symbol s cannot expand into anything that begins with t (or, if s can go to λ , nothing in follow of s leads off with t)

Grammar

► The grammar:

IF \rightarrow if\

ELSE \rightarrow else\

SEMI \rightarrow ;

NUM \rightarrow \d+

ID \rightarrow \w+

EQ \rightarrow =

LP \rightarrow \((

RP \rightarrow \)

ADDOP \rightarrow \+

MULOP \rightarrow *

LBR \rightarrow \{

RBR \rightarrow \}

S \rightarrow stmt SEMI S | λ

a-o-f \rightarrow ID a-o-f'

a-o-f' \rightarrow EQ e | LP e

RP

cond \rightarrow IF LP e RP

LBR S RBR cond'

cond' \rightarrow λ | ELSE

LBR S RBR

e \rightarrow t e'

e' \rightarrow ADDOP t e' | λ

f \rightarrow ID | NUM | LP e

RP

stmt \rightarrow a-o-f | cond

t \rightarrow f t'

t' \rightarrow MULOP f t' | λ

Table

	\$	ADDOP	ELSE	EQ	ID	IF	LP	MULOP	NUM	RBR	RP	SEMI
S	λ	•	•	•	stmt SEMI S	stmt SEMI S	•	•	•	λ	•	•
a-o-f	•	•	•	•	ID a-o-f'	•	•	•	•	•	•	•
a-o-f'	•	•	•	EQ e	•	•	LP e RP	•	•	•	•	•
cond	•	•	•	•	•	IF LP e RP LBR S RBR cond'	•	•	•	•	•	•
cond'	•	•	ELSE LBR S RBR	•	•	•	•	•	•	•	•	λ
e	•	•	•	•	t e'	•	t e'	•	t e'	•	•	•
e'	•	ADDOP t e'	•	•	•	•	•	•	•	•	λ	λ
f	•	•	•	•	ID	•	LP e RP	•	NUM	•	•	•
stmt	•	•	•	•	a-o-f	cond	•	•	•	•	•	•
t	•	•	•	•	f t'	•	f t'	•	f t'	•	•	•
t'	•	λ	•	•	•	•	•	MULOP f t'	•	•	λ	λ

Example

- ▶ Input: $x=7*4;$
- ▶ Tokens: ID, EQ, NUM, MULOP, NUM
- ▶ Trace out parse actions (in class)

Tree

- ▶ How can we get the tree?
- ▶ Suppose we have:

```
1 class TreeNode:
2     def __init__(self,sym):
3         self.children=[]
4         self.sym = sym
5         self.token = None
```

- ▶ Instead of storing strings to stack, put tree nodes on stack

Loop

```
1 while input remains:
2     let t = symbol of next token of input
3     let s = stack.top().sym
4     if s is a terminal:
5         if s == t:
6             stack.top().token = next token of input
7             consume token
8             pop stack
9     else:
10        n = stack.top()
11        pop stack
12        for sym in reversed(table[s][t]):
13            c = TreeNode(sym)
14            n.children.prepend(c)
15        push c to stack
```

Bootstrap

- ▶ To begin:
- ▶ `root = TreeNode(start_symbol)`
- ▶ `stk.push(root)`
- ▶ When done: “root” is the root of the tree

Example

- ▶ Trace previous example but this time, create tree

Assignment

- ▶ Write a program which works with the [test harness](#) to produce a correct parse tree.

Sources

- ▶ Aho, Lam, Sethi, Ullman. Compilers: Principles, Techniques, & Tools (2nd ed).
- ▶ K. Louden. Compiler Construction: Principles and Practice.

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