ASM 3

Motivation

- Work with variables
 - Also support different types
 - Primarily string and number

Grammar

- ► Change the grammar: We previously had: program → braceblock
- Now we'll have: program → var-decl-list braceblock

Grammar

- Add variable declarations: var-decl-list \rightarrow var-decl SEMI var-decl-list $\mid \lambda$ var-decl \rightarrow VAR ID type type \rightarrow non-array-type \mid non-array-type \rightarrow NUMBER \mid STRING
- Later, we might add arrays

Grammar

Finally, add assignments: stmt \rightarrow ... | assign SEMI assign \rightarrow ID EQ expr

Code

```
void assignNodeCode( TreeNode n ){
    // assign -> ID EQ expr
    VarType t0;
    exprNodeCode( n.Children[2], out t0 ); //result on stack
    emit("pop rax");
    emit("mov [{0}], rax" , n.Children[0].Token.Lexeme );
}
```

- This is our first attempt
- So, of course, it's not right. What are the ways it can fail?

Problems

- Undeclared variables
- Wrong variable type (lhs and rhs type mismatch)
- Name conflict
 - ▶ What if the user does something like: rax = 42?
 - ▶ The assembler will reject the code

Solution

- Let's tackle the last problem first!
- Recall: We have a label() function which gives a unique label
- We'll always change the names that the user gives us into our own labels
- Maintain a dictionary
 - ▶ Key = variable name as specified in user's code
 - ▶ Value = Information about that variable
 - At least the label and the type
 - We might store the source code line where it was declared for error diagnostics

Symbol Table

- What we're creating is a symbol table
- We can create a class for this
- First, a type that holds information about a single variable

Symbol Table

► The table itself:

```
class SymbolTable{
    public Dictionary<string, VarInfo> table = new Dictionary< ... >();
    public VarInfo this[string v]{
        get
            return table[v];
        set
            if( table.ContainsKev(v) )
                error: Redeclaration?
            table[v] = value:
    public bool Contains(string v){
        return table.ContainsKey(v);
```

► Then make a static: static SymbolTable symtable;

Code

 When code declares a global variable: We enter information about it into our symbol table

```
void vardeclNodeCode(TreeNode n) {
    string vname = n.Children[1].Lexeme;
    string vtypestr = n.Children[2].Children[0].Symbol;
    VarType vtype = (VarType) Enum.Parse(typeof(VarType), vtypestr
    );
    if( symtable.Contains(vname) ) {
        error: Duplicate declaration of vname
    }
    symtable[vname] = new VarInfo(vtype,label());
}
```

Assign

Now we can go back and fix assign:

```
void assignNodeCode( TreeNode n ){
    // assign -> ID EQ expr
    var a1 = exprNodeCode( n.Children[2] );
    emit("pop rax");
    string vname = n.Children[0].Lexeme;
    if( !symtable.Contains(vname) )
        error: Undeclared variable
    if( symtable[vname].VType != a1["type"] )
        error: Type mismatch
    emit("mov [{0}], rax" , symtable[vname].Label );
}
```

- Notice that we check variable type with the synthesized type attribute that came back from expr
- Last time, we set it to always use VarType.NUMBER

Analysis

- We've handled undeclared variables
- And wrong variable type (lhs and rhs type mismatch)
- And name conflicts
- Life is good!
 - ▶ Right?

Problem

- ▶ But we have a bit of a problem: We don't have any way to actually assign strings to variables
- Ex: x = "foo"
- We need to go back and tweak factor

Add two productions for factor:

```
factor → NUM | LP expr RP | STRING-CONSTANT | ID
void factorNodeCode(TreeNode n, out VarType type){
    //factor -> NUM | LP expr RP | STRING-CONSTANT | ID
    var child = n.Children[0];
    switch( child.Symbol ){
        case "NUM":
            emit("push float64 ({0})".Convert.ToDouble( child.Token.Lexeme));
            type = VarType.NUMBER:
            break:
        case "LP":
            exprNodeCode( n.Children[1], type );
            break:
        case "STRING-CONSTANT":
            ???
            type = VarType.STRING;
            break;
        case "ID":
            ???
        default:
            error
```

Decisions

- We also have to decide if we need to tweak sum and term
 - Ex: In Python, we can write things like: "abc" * 12
 - ▶ If we want to allow that, we'd add code to term
 - ▶ In languages like Java, C#, Python, we can do string+string
 - ▶ In Java and C# we can also do string+num
- For simplicity, we'll disallow both of these operations

Question

- How are strings to be represented? We can't realistically dump a bunch of text data on the stack
 - ▶ Too slow and inefficient!
- We'll define strings as immutable
- Then we'll add code for a string pool
 - Declare a global: static Dictionary<string,string> stringPool;
 - key = The string constant itself
 - ▶ value = Its label
- When we want to refer to a string, we just return its address

Note

- Observe carefully the difference:
 - ▶ mov rax, [x]
 - ▶ Moves contents of memory at label x into rax
 - mov rax, x
 - Moves the location in memory that label x refers to into rax

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- Consider C statement: "x=y;"
 - Equivalent assembly: mov rax,[y] mov [x],rax
- Consider C statement: "x=&y;"
 - Equivalent assembly: mov rax,y mov [x],rax

Problem

- Now we can write the code for the string-constant evaluation
 - ▶ The lexeme will contain the surrounding quotation marks
 - ▶ And if the string constant had backslash escapes, those will be un-expanded
- Our string-constant node will have one synthesized attribute:
 The label associated with the string constant

Now we can write factor's case for STRING-CONSTANT

```
case "STRING-CONSTANT":
    string lbl;
    stringconstantNodeCode( child, out lbl );
    //notice: No brackets!
    emit("mov rax, {0}", lbl );
    emit("push rax");
    type = VarType.STRING;
    break;
```

This stores the address of the string data on the stack

- We're ready to implement the code for the "ID" case in factor
- First, make sure the identifier exists

```
case "ID":
    string vname = n.Children[0].Lexeme;
    if( !symtable.Contains(vname) )
        error: Undeclared variable
```

- ► If the variable represents a number: The variable will contain the value itself
- If the variable represents a string: The variable contains the string's address in memory. So we just push that to the stack

```
VarInfo vi = symtable[vname];
switch(vi.VType){
   case VarType.NUMBER:
   case VarType.STRING:
       emit("mov rax,[{0}]" , symtable[vname].Label );
       emit("push rax");
       break;
   default:
       ICE
}
type = vi.VType;
```

Done Parsing

There's one more detail: When we're done parsing, we must output data from the symbol table and string pool:

```
emit("section .data");
outputSymbolTableInfo();
outputStringPoolInfo();
```

Symbol Table

String Pool

```
void outputStringPoolInfo(){
    foreach( var tmp in stringPool ){
        string the String = tmp.Key;
        string lbl = tmp.Value;
        emit( lbl+":");
        byte[] B = System.Text.Encoding.ASCII.GetBytes(theString);
        foreach( byte b in B )
            emit("db {0}", b );
        emit( "db 0" ); //null terminator
```

Assignment

- Implement the necessary code for the test harness.
 - ► Files: Main.cs, inputs.txt, GrammarData.cs

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