LR0 Table

Parsing

- ▶ It's a little awkward to use the DFA in the parse stage
- So we will build a parse table
 - ▶ Sort of like LL table
 - One row per state
 - One column per grammar symbol
- If you are in state s and next input symbol is t: Consult table[s][t] to determine what to do next
 - ▶ Recall: LR parser can either shift or reduce.

Build Table

- ► To build the table: Suppose we have a list of DFA states
- First, we assign each state a unique number:

```
for i in range(len(states)):
    states[i].index = i
```

- Or, you could have assigned these numbers when you created the DFA nodes
- Now we're ready to build the table: table = []

Build Table

```
for s in states:
    row = \{\}
    for sym in s.transitions:
        if sym is a terminal:
            row[sym] = ("S",s.transitions[sym].index)
        else:
            row[sym] = ("T",s.transitions[sym].index)
    for item in s.items:
       lhs, rhs, dpos = item
        if doos at end of rhs:
            for w in terminals:
                row[w] = ("R", len(rhs), lhs)
            row["$"] = ("R", len(rhs), lhs )
    table.append(row)l
```

Conflict

- ► If table cell set more than once: Grammar is not LR(0)
 - ▶ Shift-reduce or reduce-reduce conflict

Shift-Reduce

- ► SR conflict: Parser could shift or reduce
 - ► Remember: Shift means "consume token from input" while reduce means "convert a RHS to a LHS"
- Often, we can resolve SR conflicts by preferring shift
 - But sometimes that's not the right choice
 - Later, we'll see how we can restrict parser's choices to deal with conflicts

Reduce-Reduce

- Parser has two or more choices of items it can reduce to
- Ex: Suppose state has items:
 - $A \rightarrow B C \bullet$
 - $X \to Y \bullet$
- Should parser reduce to A or reduce to X?

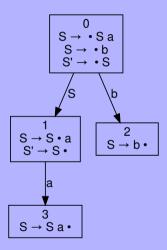
Example

Consider this grammar:

$$\begin{array}{c} a \rightarrow a \\ b \rightarrow b \end{array}$$

$$S \rightarrow S a \mid b$$

DFA



Table

	State	\$	S	a	b
0	$S \to \bullet S a$ $S \to \bullet b$		T,1		S,2
	$S' \rightarrow \bullet S$				
1	$S \rightarrow S \bullet a$ $S' \rightarrow S \bullet$	R,1,S'		S,3 R,1,S'	R,1,S'
2	$S \rightarrow b \bullet$	R,1,S		R,1,S	R,1,S
3	$S \rightarrow S a \bullet$	R,2,S		R,2,S	R,2,S

Assignment

- None!
- Next time, we'll build and use the table...

Sources

- K. Louden. Compiler Construction: Principles and Practice.
- A. Aho, M. Lam, R. Sethi, J. Ullman. Compilers: Principles, Techniques, & Tools. (2nd ed.) Addison-Wesley.

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