

# Combinatorial Optimization

Written by HADIOUCHE Azouaou.

## Disclaimer

This course extremely losely follows what Mr.Berrachedi is supposed to mean.

To seperate the contents of the course to actual additions or out of context information, a black band will be added by its side like the one englobing this comment.

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# Chapter 1

# Optimization Problems

**Note:** These will be added later... probably

- Vector comparison
- Problem notation

## 1.1. Transformation Of C.O. Problems

We consider the following integer optimization problem,

$$(INP) \quad \begin{cases} \max c \cdot x \\ Ax \leq b \\ x \in \mathbb{N}^m \end{cases}$$

with  $A \in \mathbb{Z}^{n \times m}$ ,  $b \in \mathbb{Z}^m$ ,  $c \in \mathbb{Z}^m$ .

- A special case of this is a binary variable problem (BVP), where we take the variable we want to optimize on  $x$  to be binary, that is,  $x \in \{0, 1\}^m$ . The knapsack is a special case where the  $b$  is the volume and  $A$  is the volume of each element, while  $c$  is their usefulness and  $x$  represents either taking it or not.

If we take by hypothesis that  $\mathcal{D}$  is bounded, notice that the set  $S \subset \mathcal{D}$  that satisfies the INP problem are finite.

**Note:** Proof will be written in a bit