Stanford University ICPC Team Notebook (2015-16)

Contents

1 Combinatorial optimization

1.1 Sparse max-flow

```
// Adjacency list implementation of Dinic's blocking flow algorithm.
// This is very fast in practice, and only loses to push-relabel flow.
// Running time:
      O(|V|^2 |E|)
       - graph, constructed using AddEdge()
       - source and sink
// OUTPUT:
       - maximum flow value
       - To obtain actual flow values, look at edges with capacity > 0
         (zero capacity edges are residual edges).
#include<cstdio>
#include < vector >
#include < queue >
using namespace std;
typedef long long LL;
struct Edge {
  int u, v;
  LL cap, flow;
Edge() {}
  Edge(int u, int v, LL cap): u(u), v(v), cap(cap), flow(0) {}
struct Dinic {
  int N;
  vector<Edge> E;
  vector<vector<int>> g;
  vector<int> d, pt;
  Dinic(int N): N(N), E(0), g(N), d(N), pt(N) {}
  void AddEdge(int u, int v, LL cap) {
    if (u != v) {
      E.emplace_back(u, v, cap);
      g[u].emplace_back(E.size() - 1);
E.emplace_back(v, u, 0);
      g[v].emplace_back(E.size() - 1);
  bool BFS(int S, int T) {
    queue<int> q({S});
     fill(d.begin(), d.end(), N + 1);
    while(!q.empty()) {
  int u = q.front(); q.pop();
  if (u == T) break;
      for (int k: g[u]) {
        Edge &e = E[k];
        if (e.flow < e.cap && d[e.v] > d[e.u] + 1) {
          d[e.v] = d[e.u] + 1;
           q.emplace(e.v);
    return d[T] != N + 1;
  LL DFS(int u, int T, LL flow = -1) {
    if (u == T || flow == 0) return flow;
    for (int &i = pt[u]; i < g[u].size(); ++i) {</pre>
      Edge &e = E[g[u][i]];
Edge &oe = E[g[u][i]^1];
if (d[e.v] == d[e.u] + 1) {
        LL amt = e.cap - e.flow;
        if (flow != -1 && amt > flow) amt = flow;
        if (LL pushed = DFS(e.v, T, amt)) {
```

```
e.flow += pushed;
          oe.flow -= pushed;
          return pushed;
    return 0;
  LL MaxFlow(int S, int T) {
    LL total = 0;
    while (BFS(S, T)) {
     fill(pt.begin(), pt.end(), 0);
while (LL flow = DFS(S, T))
       total += flow;
    return total:
// BEGIN CUT
// The following code solves SPOJ problem #4110: Fast Maximum Flow (FASTFLOW)
  int N. E:
  scanf("%d%d", &N, &E);
  Dinic dinic(N);
  for (int i = 0; i < E; i++)
    int u. v:
    LL cap;
    scanf("%d%d%lld", &u, &v, &cap);
    dinic.AddEdge(u - 1, v - 1, cap);
    dinic.AddEdge(v - 1, u - 1, cap);
  printf("%1ld\n", dinic.MaxFlow(0, N - 1));
  return 0;
// END CUT
```

1.2 Min-cost max-flow

```
// Implementation of min cost max flow algorithm using adjacency
// matrix (Edmonds and Karp 1972). This implementation keeps track of
// forward and reverse edges separately (so you can set cap[i][j] !=
// cap[j][i]). For a regular max flow, set all edge costs to 0.
// Running time, O(|V|^2) cost per augmentation
       max flow:
                           O(|V|^3) augmentations
       min cost max flow: O(|V|^4 * MAX\_EDGE\_COST) augmentations
// INPUT:
       - graph, constructed using AddEdge()
       - source
      - sink
       - (maximum flow value, minimum cost value)
       - To obtain the actual flow, look at positive values only.
#include <bits/stdc++.h>
using namespace std;
typedef vector<int> VI;
typedef vector<VI> VVI;
typedef long long L;
typedef vector<L> VL;
typedef vector<VL> VVL;
typedef pair<int, int> PII;
typedef vector<PII> VPII;
const L INF = numeric_limits<L>::max() / 4;
struct MinCostMaxFlow {
    int N;
    VVL cap, flow, cost;
    VI found;
    VL dist, pi, width;
    VPII dad;
    MinCostMaxFlow(int N) :
        N(N), cap(N, VL(N)), flow(N, VL(N)), cost(N, VL(N)), found(N), dist(N), pi(N), width(N), dad(N) {}
    void AddEdge(int from, int to, L cap, L cost) {
```

```
this->cap[from][to] = cap;
         this->cost[from][to] = cost;
    void Relax(int s, int k, L cap, L cost, int dir) {
         L \text{ val} = \text{dist}[s] + \text{pi}[s] - \text{pi}[k] + \text{cost};
         if (cap && val < dist[k]) {</pre>
             dist[k] = val;
             dad[k] = make_pair(s, dir);
width[k] = min(cap, width[s]);
    L Dijkstra(int s, int t) {
         fill(found.begin(), found.end(), false);
fill(dist.begin(), dist.end(), INF);
         fill(width.begin(), width.end(), 0);
         dist[s] = 0;
         width[s] = INF;
         while (s != -1) {
             int best = -1;
              found[s] = true;
             for (int k = 0; k < N; k++) {
                 if (found[k]) continue;
                 Relax(s, k, cap[s][k] - flow[s][k], cost[s][k], 1);
Relax(s, k, flow[k][s], -cost[k][s], -1);
                 if (best == -1 || dist[k] < dist[best]) best = k;</pre>
             s = best;
         for (int k = 0; k < N; k++)
             pi[k] = min(pi[k] + dist[k], INF);
         return width[t];
    pair<L, L> GetMaxFlow(int s, int t) {
         L totflow = 0, totcost = 0;
         while (L amt = Dijkstra(s, t)) {
   totflow += amt;
             for (int x = t; x != s; x = dad[x].first) {
                 if (dad[x].second == 1) {
                      flow[dad[x].first][x] += amt;
                       totcost += amt * cost[dad[x].first][x];
                  } else {
                      flow[x][dad[x].first] -= amt;
                      totcost -= amt * cost[x][dad[x].first];
         return make_pair(totflow, totcost);
1:
// BEGIN CUT
// The following code solves UVA problem #10594: Data Flow
int main() {
    int N, M;
    while (scanf("%d%d", &N, &M) == 2) {
         VVL v(M, VL(3));
         for (int i = 0; i < M; i++)
             scanf("%Ld%Ld%Ld", &v[i][0], &v[i][1], &v[i][2]);
         L D, K;
         scanf("%Ld%Ld", &D, &K);
         MinCostMaxFlow mcmf(N+1);
        for (int i = 0; i < M; i++) {
    mcmf.AddEdge(int(v[i][0]), int(v[i][1]), K, v[i][2]);</pre>
             mcmf.AddEdge(int(v[i][1]), int(v[i][0]), K, v[i][2]);
         mcmf.AddEdge(0, 1, D, 0);
         pair<L, L> res = mcmf.GetMaxFlow(0, N);
         if (res.first == D) {
             printf("%Ld\n", res.second);
         } else {
             printf("Impossible.\n");
    return 0:
// END CUT
```

1.3 Push-relabel max-flow

```
// Adjacency list implementation of FIFO push relabel maximum flow
// with the gap relabeling heuristic. This implementation is
// significantly faster than straight Ford-Fulkerson. It solves
// random problems with 10000 vertices and 1000000 edges in a few
// seconds, though it is possible to construct test cases that
// achieve the worst-case.
// Running time:
      0(|V|^3)
// INPUT:
       - graph, constructed using AddEdge()
       - source
      - sink
// OUTPUT:
       - maximum flow value
       - To obtain the actual flow values, look at all edges with
         capacity > 0 (zero capacity edges are residual edges).
#include <cmath>
#include <vector>
#include <iostream>
#include <queue>
using namespace std:
typedef long long LL;
 int from, to, cap, flow, index;
Edge(int from, int to, int cap, int flow, int index) :
    from(from), to(to), cap(cap), flow(flow), index(index) {}
};
struct PushRelabel {
 int N;
  vector<vector<Edge> > G:
  vector<LL> excess:
  vector<int> dist, active, count;
  queue<int> Q;
  PushRelabel(int N) : N(N), G(N), excess(N), dist(N), active(N), count(2*N) {}
  void AddEdge(int from, int to, int cap)
    G[from].push_back(Edge(from, to, cap, 0, G[to].size()));
    if (from == to) G[from].back().index++;
    G[to].push\_back(Edge(to, from, 0, 0, G[from].size() - 1));
  void Enqueue(int v) {
   if (!active[v] && excess[v] > 0) { active[v] = true; Q.push(v); }
  void Push (Edge &e) {
    int amt = int(min(excess[e.from], LL(e.cap - e.flow)));
    if (dist[e.from] <= dist[e.to] || amt == 0) return;</pre>
    e.flow += amt;
    G[e.to][e.index].flow -= amt;
    excess[e.to] += amt;
    excess[e.from] -= amt;
    Enqueue (e.to);
  void Gap(int k) {
    for (int v = 0; v < N; v++) {
      if (dist[v] < k) continue;</pre>
      count[dist[v]]--:
      dist[v] = max(dist[v], N+1);
      count[dist[v]]++;
      Enqueue (v):
  void Relabel(int v) {
    count [dist[v]]--;
    dist[v] = 2*N;
    for (int i = 0; i < G[v].size(); i++)</pre>
     if (G[v][i].cap - G[v][i].flow > 0)
        dist[v] = min(dist[v], dist[G[v][i].to] + 1);
    count[dist[v]]++;
    Enqueue (v);
  void Discharge(int v) {
    for (int i = 0; excess[v] > 0 && i < G[v].size(); i++) Push(G[v][i]);</pre>
    if (excess[v] > 0) {
```

2

```
if (count[dist[v]] == 1)
         Gap(dist[v]);
      else
        Relabel(v);
  LL GetMaxFlow(int s, int t) {
    count[0] = N-1;
count[N] = 1;
   dist[s] = N;
active[s] = active[t] = true;
for (int i = 0; i < G[s].size(); i++) {
   excess[s] += G[s][i].cap;</pre>
      Push(G[s][i]);
    while (!Q.empty()) {
      int v = Q.front();
       active[v] = false;
      Discharge(v);
    return totflow:
1:
// The following code solves SPOJ problem #4110: Fast Maximum Flow (FASTFLOW)
int main() {
  scanf("%d%d", &n, &m);
  PushRelabel pr(n);
  for (int i = 0; i < m; i++) {
   int a, b, c;
scanf("%d%d%d", &a, &b, &c);
    if (a == b) continue;
    pr.AddEdge(a-1, b-1, c);
pr.AddEdge(b-1, a-1, c);
  printf("%Ld\n", pr.GetMaxFlow(0, n-1));
  return 0;
// END CUT
```

1.4 Min-cost matching

```
// Min cost bipartite matching via shortest augmenting paths
// This is an O(n^3) implementation of a shortest augmenting path
// algorithm for finding min cost perfect matchings in dense
// graphs. In practice, it solves 1000x1000 problems in around 1
// second.
    cost[i][j] = cost for pairing left node i with right node j
    Lmate[i] = index of right node that left node i pairs with
    Rmate[j] = index of left node that right node j pairs with
// The values in cost[i][j] may be positive or negative. To perform
// maximization, simply negate the cost[][] matrix.
#include <algorithm>
#include <cstdio>
#include <cmath>
#include <vector>
using namespace std;
typedef vector<double> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
double MinCostMatching(const VVD &cost, VI &Lmate, VI &Rmate) {
  int n = int(cost.size());
  // construct dual feasible solution
  VD u(n);
  VD v(n);
  for (int i = 0; i < n; i++) {</pre>
```

```
u[i] = cost[i][0];
  for (int j = 1; j < n; j++) u[i] = min(u[i], cost[i][j]);</pre>
for (int j = 0; j < n; j++) {
  v[j] = cost[0][j] - u[0];
  for (int i = 1; i < n; i++) v[j] = min(v[j], cost[i][j] - u[i]);</pre>
// construct primal solution satisfying complementary slackness
Lmate = VI(n, -1);
Rmate = VI(n, -1);
Lmate[i] = j;
      Rmate[j] = i;
      mated++;
      break;
VD dist(n);
VT dad(n):
VI seen(n):
// repeat until primal solution is feasible
while (mated < n) {
   // find an unmatched left node
  while (Lmate[s] != -1) s++;
  // initialize Dijkstra
  fill(dad.begin(), dad.end(), -1);
  fill(seen.begin(), seen.end(), 0);
  for (int k = 0; k < n; k++)
    dist[k] = cost[s][k] - u[s] - v[k];
  int. i = 0:
  while (true) {
    // find closest
     i = -1;
    for (int k = 0; k < n; k++) {
      if (seen[k]) continue;
      if (j == -1 || dist[k] < dist[j]) j = k;</pre>
    seen[j] = 1;
    // termination condition
if (Rmate[j] == -1) break;
    // relax neighbors
    const int i = Rmate[j];
    for (int k = 0; k < n; k++) {
      if (seen[k]) continue;
      const double new_dist = dist[j] + cost[i][k] - u[i] - v[k];
      if (dist[k] > new_dist) {
        dist[k] = new_dist;
         dad[k] = j;
  // update dual variables
  for (int k = 0; k < n; k++) {
   if (k == j || !seen[k]) continue;
   const int i = Rmate[k];</pre>
    v[k] += dist[k] - dist[j];
    u[i] -= dist[k] - dist[j];
  u[s] += dist[j];
   // augment along path
  while (dad[j] >= 0) {
    const int d = dad[j];
    Rmate[j] = Rmate[d];
    Lmate[Rmate[j]] = j;
    j = d;
  Rmate[j] = s;
  Lmate[s] = j;
  mated++;
double value = 0;
for (int i = 0; i < n; i++)
  value += cost[i][Lmate[i]];
```

1.5 Max bipartite matchine

```
// This code performs maximum bipartite matching.
// Running time: O(|E| |V|) -- often much faster in practice
     INPUT: w[i][j] = edge between row node i and column node j
     OUTPUT: mr[i] = assignment for row node i, -1 if unassigned
             mc[j] = assignment for column node j, -1 if unassigned
              function returns number of matches made
#include <vector>
using namespace std;
typedef vector<int> VI;
typedef vector<VI> VVI;
bool FindMatch(int i, const VVI &w, VI &mr, VI &mc, VI &seen) {
  for (int j = 0; j < w[i].size(); j++) {
  if (w[i][j] && !seen[j]) {</pre>
      seen[j] = true;
      if (mc[j] < 0 || FindMatch(mc[j], w, mr, mc, seen)) {</pre>
       mr[i] = j;
        mc[j] = i;
        return true;
  return false:
int BipartiteMatching(const VVI &w, VI &mr, VI &mc) {
 mr = VI(w.size(), -1);
  mc = VI(w[0].size(), -1);
  for (int i = 0; i < w.size(); i++) {</pre>
    VI seen(w[0].size());
    if (FindMatch(i, w, mr, mc, seen)) ct++;
  return ct;
```

1.6 Global min-cut

```
// Adjacency matrix implementation of Stoer-Wagner min cut algorithm.
// Running time:
      0(|V|^3)
      - graph, constructed using AddEdge()
      - (min cut value, nodes in half of min cut)
#include <cmath>
#include <vector>
#include <iostream>
using namespace std:
typedef vector<int> VI;
typedef vector<VI> VVI;
const int INF = 1000000000;
pair<int, VI> GetMinCut(VVI &weights) {
  int N = weights.size();
  VI used(N), cut, best_cut;
  int best_weight = -1;
  for (int phase = N-1; phase >= 0; phase--) {
   VI w = weights[0];
    VI added = used;
    int prev, last = 0;
    for (int i = 0; i < phase; i++) {</pre>
     prev = last;
```

```
last = -1;
       for (int j = 1; j < N; j++)
         if (!added[j] && (last == -1 || w[j] > w[last])) last = j;
       if (i == phase-1) {
        for (int j = 0; j < N; j++) weights[prev][j] += weights[last][j];
for (int j = 0; j < N; j++) weights[j][prev] = weights[prev][j];</pre>
         used[last] = true;
         cut.push_back(last);
         if (best_weight == -1 || w[last] < best_weight) {</pre>
           best_cut = cut;
           best_weight = w[last];
       } else {
         for (int j = 0; j < N; j++)
           w[j] += weights[last][j];
         added[last] = true;
  return make_pair(best_weight, best_cut);
// BEGIN CUT
// The following code solves UVA problem #10989: Bomb, Divide and Conquer
int main() {
 int N:
  cin >> N:
  for (int i = 0; i < N; i++) {
    int n. m:
    cin >> n >> m;
     VVI weights(n, VI(n));
    for (int j = 0; j < m; j++) {
      int a, b, c;
      cin >> a >> b >> c;
      weights[a-1][b-1] = weights[b-1][a-1] = c;
    pair<int, VI> res = GetMinCut(weights);
cout << "Case #" << i+1 << ": " << res.first << endl;</pre>
// END CUT
```

1.7 Graph cut inference

```
// Special-purpose {0,1} combinatorial optimization solver for
// problems of the following by a reduction to graph cuts:
                          sum_i psi_i(x[i])
// x[1]...x[n] in {0,1} + sum_{i < j} phi_{ij}(x[i], x[j])
// where
       psi_i : {0, 1} --> R
    phi_{ij} : {0, 1} x {0, 1} --> R
// such that
    phi_{ij}(0,0) + phi_{ij}(1,1) \le phi_{ij}(0,1) + phi_{ij}(1,0)  (*)
// This can also be used to solve maximization problems where the
// direction of the inequality in (*) is reversed.
// INPUT: phi -- a \ matrix \ such \ that \ phi[i][j][u][v] = phi_{ij}(u, \ v)
         psi -- a matrix such that psi[i][u] = psi_i(u)
          x -- a vector where the optimal solution will be stored
// OUTPUT: value of the optimal solution
// To use this code, create a GraphCutInference object, and call the
// DoInference() method. To perform maximization instead of minimization,
// ensure that #define MAXIMIZATION is enabled.
#include <vector>
#include <iostream>
using namespace std;
typedef vector<int> VI;
typedef vector<VI> VVI;
typedef vector<VVI> VVVI;
typedef vector<VVVI> VVVVI;
const int INF = 10000000000:
// comment out following line for minimization
#define MAXIMIZATION
struct GraphCutInference {
 int N;
```

```
VVI cap, flow;
  VI reached;
  int Augment(int s, int t, int a) {
    reached[s] = 1;
    if (s == t) return a;
    for (int k = 0; k < N; k++) {
      if (reached[k]) continue;
       \textbf{if (int } aa = min(a, \ cap[s][k] - flow[s][k])) \ \{ \\
        if (int b = Augment(k, t, aa)) {
          flow[s][k] += b;
          flow[k][s] -= b;
          return b;
    return 0:
  int GetMaxFlow(int s, int t) {
    N = cap.size();
    flow = VVI(N, VI(N));
    reached = VI(N);
   int totflow = 0;
   while (int amt = Augment(s, t, INF)) {
      totflow += amt:
      fill(reached.begin(), reached.end(), 0);
    return totflow:
  int DoInference(const VVVVI &phi, const VVI &psi, VI &x) {
    int M = phi.size();
    cap = VVI(M+2, VI(M+2));
    VI b(M);
    int c = 0;
    for (int i = 0; i < M; i++) {
     b[i] += psi[i][1] - psi[i][0];
      c += psi[i][0];
     c += phi[i][j][0][0];
#ifdef MAXIMIZATION
   for (int i = 0; i < M; i++) {
  for (int j = i+1; j < M; j++)
    cap[i][j] *= -1;</pre>
     b[i] *= -1;
    c *= -1:
#endif
    for (int i = 0; i < M; i++) {</pre>
     if (b[i] >= 0) {
       cap[M][i] = b[i];
       cap[i][M+1] = -b[i];
        c += b[i];
    int score = GetMaxFlow(M, M+1);
   fill(reached.begin(), reached.end(), 0);
    Augment (M, M+1, INF);
    x = VI(M);
    for (int i = 0; i < M; i++) x[i] = reached[i] ? 0 : 1;</pre>
#ifdef MAXIMIZATION
    score *= -1;
#endif
    return score;
};
int main() {
  // solver for "Cat vs. Dog" from NWERC 2008
  int numcases;
  cin >> numcases;
  for (int caseno = 0; caseno < numcases; caseno++) {</pre>
   int c, d, v;
    cin >> c >> d >> v;
```

```
VVVVI phi(c+d, VVVI(c+d, VVI(2, VI(2))));
  VVI psi(c+d, VI(2));
  for (int i = 0; i < v; i++) {
    char p, q;
    int u, v;
   cin >> p >> u >> q >> v;
    u--; v--;
   if (p == 'C') {
     phi[u][c+v][0][0]++;
      phi[c+v][u][0][0]++;
    } else {
      phi[v][c+u][1][1]++;
      phi[c+u][v][1][1]++;
  GraphCutInference graph;
  cout << graph.DoInference(phi, psi, x) << endl;</pre>
return 0;
```

2 Geometry

2.1 Convex hull

```
// Compute the 2D convex hull of a set of points using the monotone chain
// algorithm. Eliminate redundant points from the hull if REMOVE_REDUNDANT is
// #defined.
// Running time: O(n log n)
    INPUT: a vector of input points, unordered.
    OUTPUT: a vector of points in the convex hull, counterclockwise, starting
             with bottommost/leftmost point
#include <cstdio>
#include <cassert>
#include <vector>
#include <algorithm>
#include <cmath>
#include <map>
using namespace std;
#define REMOVE REDUNDANT
typedef double T:
const T EPS = 1e-7;
struct PT {
  T x, y;
  PT() {}
  PT(T x, T y) : x(x), y(y) {}
  bool operator<(const PT &rhs) const { return make_pair(y,x) < make_pair(rhs.y,rhs.x); }</pre>
 bool operator == (const PT &rhs) const { return make_pair(y,x) == make_pair(rhs.y,rhs.x); }
T cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
T area2(PT a, PT b, PT c) { return cross(a,b) + cross(b,c) + cross(c,a); }
#ifdef REMOVE REDUNDANT
bool between (const PT &a, const PT &b, const PT &c) {
 return (fabs(area2(a,b,c)) < EPS && (a.x-b.x)*(c.x-b.x) <= 0 && (a.y-b.y)*(c.y-b.y) <= 0);
void ConvexHull(vector<PT> &pts) {
  sort(pts.begin(), pts.end());
  pts.erase(unique(pts.begin(), pts.end()), pts.end());
  vector<PT> up, dn;
  for (int i = 0; i < pts.size(); i++) {</pre>
    while (up.size() > 1 && area2(up[up.size()-2], up.back(), pts[i]) >= 0) up.pop_back();
    while (dn.size() > 1 && area2(dn[dn.size()-2], dn.back(), pts[i]) <= 0) dn.pop_back();</pre>
    up.push_back(pts[i]);
    dn.push_back(pts[i]);
  pts = dn;
  for (int i = (int) up.size() - 2; i >= 1; i--) pts.push back(up[i]);
#ifdef REMOVE_REDUNDANT
```

```
if (pts.size() <= 2) return;</pre>
  dn.clear();
  dn.push_back(pts[0]);
  dn.push_back(pts[1]);
  for (int i = 2; i < pts.size(); i++) {</pre>
    if (between(dn[dn.size()-2], dn[dn.size()-1], pts[i])) dn.pop_back();
    dn.push_back(pts[i]);
  if (dn.size() >= 3 && between(dn.back(), dn[0], dn[1])) {
    dn[0] = dn.back();
    dn.pop_back();
  pts = dn;
#endif
// The following code solves SPOJ problem #26: Build the Fence (BSHEEP)
int main()
 int t;
  scanf("%d", &t);
  for (int caseno = 0; caseno < t; caseno++) {</pre>
    int n:
    scanf("%d", &n);
    vector<PT> v(n);
    for (int i = 0; i < n; i++) scanf("%lf%lf", &v[i].x, &v[i].y);</pre>
    vector<PT> h(v):
    map<PT.int> index:
    for (int i = n-1; i >= 0; i--) index[v[i]] = i+1;
    ConvexHull(h);
    for (int i = 0; i < h.size(); i++) {</pre>
      double dx = h[i].x - h[(i+1)%h.size()].x;
      double dy = h[i].y - h[(i+1)%h.size()].y;
      len += sqrt (dx*dx+dy*dy);
    if (caseno > 0) printf("\n");
printf("%.2f\n", len);
    for (int i = 0; i < h.size(); i++) {
  if (i > 0) printf(" ");
      printf("%d", index[h[i]]);
    printf("\n");
// END CUT
```

2.2 Miscellaneous geometry

```
// C++ routines for computational geometry.
#include <iostream>
#include <vector>
#include <cmath>
#include <cassert>
using namespace std;
double INF = 1e100:
double EPS = 1e-12;
struct PT (
  double x, y;
  PT() {}
  PT (double x, double y) : x(x), y(y) {}
  PT (const PT &p) : x(p.x), y(p.y)
  PT operator + (const PT &p) const { return PT(x+p.x, y+p.y); }
  PT operator - (const PT &p) const { return PT(x-p.x, y-p.y); }
                                 const { return PT(x*c, y*c ); ]
  PT operator * (double c)
  PT operator / (double c)
                                const { return PT(x/c, y/c ); }
                              { return p.x*q.x+p.y*q.y; }
double dist2(PT p, PT q) { return dot(p-q,p-q); }
double cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
ostream &operator<<(ostream &os, const PT &p) {
    return os << "(" << p.x << "," << p.y << ")";
// rotate a point CCW or CW around the origin
PT RotateCCW90 (PT p) { return PT(-p.y,p.x); }
PT RotateCW90 (PT p) { return PT(p.y,-p.x); }
                         { return PT(p.y,-p.x); }
PT RotateCCW(PT p, double t) {
```

```
return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos(t));
// project point c onto line through a and b
PT ProjectPointLine(PT a, PT b, PT c) {
  return a + (b-a) *dot (c-a, b-a) /dot (b-a, b-a);
// project point c onto line segment through a and b
PT ProjectPointSegment(PT a, PT b, PT c) {
  double r = dot(b-a,b-a);
  if (fabs(r) < EPS) return a;</pre>
  r = dot(c-a, b-a)/r;
  if (r < 0) return a:
  if (r > 1) return b;
 return a + (b-a) *r;
// compute distance from c to segment between a and b
double DistancePointSegment (PT a, PT b, PT c) {
 return sqrt(dist2(c, ProjectPointSegment(a, b, c)));
// compute distance between point (x,y,z) and plane ax+by+cz=d
double DistancePointPlane(double x, double y, double z,
                           double a, double b, double c, double d)
  return fabs(a*x+b*v+c*z-d)/sgrt(a*a+b*b+c*c);
// determine if lines from a to b and c to d are parallel or collinear
bool LinesParallel(PT a, PT b, PT c, PT d) {
 return fabs(cross(b-a, c-d)) < EPS;
bool LinesCollinear(PT a, PT b, PT c, PT d) {
  return LinesParallel(a, b, c, d)
      && fabs(cross(a-b, a-c)) < EPS
      && fabs(cross(c-d, c-a)) < EPS;
// determine if line segment from a to b intersects with
// line segment from c to d
bool SegmentsIntersect(PT a, PT b, PT c, PT d) {
 if (LinesCollinear(a, b, c, d)) {
   if (dist2(a, c) < EPS || dist2(a, d) < EPS ||</pre>
      dist2(b, c) < EPS || dist2(b, d) < EPS) return true;
    if (dot(c-a, c-b) > 0 && dot(d-a, d-b) > 0 && dot(c-b, d-b) > 0)
     return false:
    return true:
  if (cross(d-a, b-a) * cross(c-a, b-a) > 0) return false;
  if (cross(a-c, d-c) * cross(b-c, d-c) > 0) return false;
  return true:
// compute intersection of line passing through a and b
// with line passing through c and d, assuming that unique
// intersection exists; for segment intersection, check if
// segments intersect first
PT ComputeLineIntersection(PT a, PT b, PT c, PT d) {
 b=b-a; d=c-d; c=c-a;
  assert(dot(b, b) > EPS && dot(d, d) > EPS);
  return a + b*cross(c, d)/cross(b, d);
// compute center of circle given three points
PT ComputeCircleCenter(PT a, PT b, PT c) {
 b = (a+b)/2:
 c = (a+c)/2:
 return ComputeLineIntersection(b, b+RotateCW90(a-b), c, c+RotateCW90(a-c));
// determine if point is in a possibly non-convex polygon (by William
// Randolph Franklin); returns 1 for strictly interior points, 0 for
// strictly exterior points, and 0 or 1 for the remaining points.
// Note that it is possible to convert this into an *exact* test using
// integer arithmetic by taking care of the division appropriately
// (making sure to deal with signs properly) and then by writing exact
// tests for checking point on polygon boundary
bool PointInPolygon(const vector<PT> &p, PT q) {
  bool c = 0:
  for (int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) %p.size();
   iff ([p[i].y <= q.y && q.y < p[j].y ||
    p[j].y <= q.y && q.y < p[i].y) &&
    q.x < p[i].x + (p[j].x - p[i].x) * (q.y - p[i].y) / (p[j].y - p[i].y))</pre>
  return c:
```

```
// determine if point is on the boundary of a polygon
bool PointOnPolygon(const vector<PT> &p, PT q) {
  for (int i = 0; i < p.size(); i++)</pre>
    if (dist2(ProjectPointSegment(p[i], p[(i+1)%p.size()], q), q) < EPS)</pre>
      return true;
    return false;
// compute intersection of line through points a and b with
// circle centered at c with radius r > 0
vector<PT> CircleLineIntersection(PT a, PT b, PT c, double r) {
  vector<PT> ret;
  b = b-a:
  a = a-c:
  double A = dot(b, b);
  double B = dot(a, b);
  double C = dot(a, a) - r*r;
  double D = B*B - A*C;
  if (D < -EPS) return ret;
  ret.push_back(c+a+b*(-B+sqrt(D+EPS))/A);
  if (D > EPS)
    ret.push_back(c+a+b*(-B-sqrt(D))/A);
  return ret:
// compute intersection of circle centered at a with radius r
// with circle centered at b with radius R
vector<PT> CircleCircleIntersection(PT a, PT b, double r, double R) {
  vector<PT> ret:
  double d = sqrt(dist2(a, b));
  if (d > r+R | | d+min(r, R) < max(r, R)) return ret;</pre>
  double x = (d*d-R*R+r*r)/(2*d);
  double y = sqrt(r*r-x*x);
  PT v = (b-a)/d;
  ret.push_back(a+v*x + RotateCCW90(v)*y);
  if (y > 0)
    ret.push_back(a+v*x - RotateCCW90(v)*y);
  return ret;
// This code computes the area or centroid of a (possibly nonconvex)
// polygon, assuming that the coordinates are listed in a clockwise or
// counterclockwise fashion. Note that the centroid is often known as // the "center of gravity" or "center of mass".
double ComputeSignedArea(const vector<PT> &p) {
  double area = 0;
  for(int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
    area += p[i].x*p[j].y - p[j].x*p[i].y;
  return area / 2.0;
double ComputeArea(const vector<PT> &p) {
 return fabs (ComputeSignedArea(p));
PT ComputeCentroid(const vector<PT> &p) {
  PT c(0,0);
  double scale = 6.0 * ComputeSignedArea(p);
  for (int i = 0; i < p.size(); i++) {
    int j = (i+1) % p.size();
    c = c + (p[i]+p[j])*(p[i].x*p[j].y - p[j].x*p[i].y);
  return c / scale;
// tests whether or not a given polygon (in CW or CCW order) is simple \,
bool IsSimple(const vector<PT> &p)
  for (int i = 0; i < p.size(); i++) {
  for (int k = i+1; k < p.size(); k++) {</pre>
     int j = (i+1) % p.size();
int l = (k+1) % p.size();
      if (i == 1 \mid \mid j == k) continue;
      if (SegmentsIntersect(p[i], p[j], p[k], p[l]))
  return true;
int main() {
  // expected: (-5,2)
  cerr << RotateCCW90(PT(2,5)) << endl;</pre>
  // expected: (5,-2)
  cerr << RotateCW90(PT(2,5)) << endl;</pre>
  // expected: (-5,2)
  cerr << RotateCCW(PT(2,5),M_PI/2) << endl;
```

```
// expected: (5,2)
cerr << ProjectPointLine(PT(-5,-2), PT(10,4), PT(3,7)) << endl;</pre>
// expected: (5,2) (7.5,3) (2.5,1)
cerr << ProjectPointSegment(PT(-5,-2), PT(10,4), PT(3,7)) << " "</pre>
     << ProjectPointSegment(PT(7.5,3), PT(10,4), PT(3,7)) << " "
     << ProjectPointSegment(PT(-5,-2), PT(2.5,1), PT(3,7)) << endl;
// expected: 6.78903
cerr << DistancePointPlane(4,-4,3,2,-2,5,-8) << endl;</pre>
// expected: 1 0 1
<< LinesParallel(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;
cerr << LinesCollinear(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "
     << LinesCollinear(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << " "
     << LinesCollinear(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;
// expected: 1 1 1 0
cerr << SegmentsIntersect(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << " "
     << SegmentsIntersect(PT(0,0), PT(2,4), PT(4,3), PT(0,5)) << " "
     << SegmentsIntersect(PT(0,0), PT(2,4), PT(2,-1), PT(-2,1)) << " "
     << SegmentsIntersect(PT(0,0), PT(2,4), PT(5,5), PT(1,7)) << endl;
// expected: (1.2)
cerr << ComputeLineIntersection(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << endl;</pre>
// expected: (1.1)
cerr << ComputeCircleCenter(PT(-3,4), PT(6,1), PT(4,5)) << endl;</pre>
v.push_back(PT(0,0));
v.push_back(PT(5,0));
v.push back(PT(5,5));
v.push_back(PT(0,5));
// expected: 1 1 1 0 0
cerr << PointInPolygon(v, PT(2,2)) << " "</pre>
     << PointInPolygon(v, PT(2,0)) << " "
     << PointInPolygon(v, PT(0,2)) << " "
<< PointInPolygon(v, PT(5,2)) << " "</pre>
     << PointInPolygon(v, PT(2,5)) << endl;
// expected: 0 1 1 1 1
cerr << PointOnPolygon(v, PT(2,2)) << " "
     << PointOnPolygon(v, PT(2,0)) << " "
     << PointOnPolygon(v, PT(0,2)) << " "</pre>
     << PointOnPolygon(v, PT(5,2)) << " "
     << PointOnPolygon(v, PT(2,5)) << endl;
// expected: (1.6)
              (5,4) (4,5)
              blank line
              (4.5) (5.4)
              blank line
              (4.5) (5.4)
vector<PT> u = CircleLineIntersection(PT(0,6), PT(2,6), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
  = CircleLineIntersection(PT(0,9), PT(9,0), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
u = CircleCircleIntersection(PT(1,1), PT(10,10), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
 u = CircleCircleIntersection (PT(1,1), \ PT(8,8), \ 5, \ 5); \\  for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl; 
u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 10, sqrt(2.0)/2.0);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 5, sqrt(2.0)/2.0);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
// area should be 5.0
// centroid should be (1.1666666, 1.166666)
PT pa[] = { PT(0,0), PT(5,0), PT(1,1), PT(0,5) };
vector<PT> p(pa, pa+4);
PT c = ComputeCentroid(p);
cerr << "Area: " << ComputeArea(p) << endl;
cerr << "Centroid: " << c << endl;
return 0;
```

2.3 Slow Delaunay triangulation

```
// degenerate cases (from O'Rourke, Computational Geometry in C)
// Running time: O(n^4)
             x[] = x-coordinates
             y[] = y-coordinates
// OUTPUT: triples = a vector containing m triples of indices
                        corresponding to triangle vertices
#include<vector>
using namespace std;
typedef double T:
struct triple {
    int i, j, k;
    triple() {}
    triple(int i, int j, int k) : i(i), j(j), k(k) {}
vector<triple> delaunayTriangulation(vector<T>& x, vector<T>& y) {
        int n = x.size();
        vector<T> z(n):
        vector<triple> ret;
        for (int i = 0; i < n; i++)
   z[i] = x[i] * x[i] + y[i] * y[i];</pre>
        for (int i = 0; i < n-2; i++) {
            for (int j = i+1; j < n; j++) {
   for (int k = i+1; k < n; k++) {</pre>
                     if (j == k) continue;
                     double xn = (y[j]-y[i])*(z[k]-z[i]) - (y[k]-y[i])*(z[j]-z[i]);
                     double yn = (x[k]-x[i])*(z[j]-z[i]) - (x[j]-x[i])*(z[k]-z[i]);
                     double zn = (x[j]-x[i]) * (y[k]-y[i]) - (x[k]-x[i]) * (y[j]-y[i]);
                     bool flag = zn < 0;</pre>
                     for (int m = 0; flag && m < n; m++)</pre>
                          flag = flag && ((x[m]-x[i])*xn +
                                           (y[m]-y[i])*yn +
                                           (z[m]-z[i])*zn <= 0);
                     if (flag) ret.push_back(triple(i, j, k));
        return ret;
int main()
    T xs[]={0, 0, 1, 0.9};
    T ys[]=\{0, 1, 0, 0.9\};
    vector<T> x(&xs[0], &xs[4]), y(&ys[0], &ys[4]);
    vector<triple> tri = delaunayTriangulation(x, y);
    //expected: 0 1 3
              0 3 2
    int i:
    for(i = 0; i < tri.size(); i++)</pre>
        printf("%d %d %d\n", tri[i].i, tri[i].j, tri[i].k);
```

3 Numerical algorithms

3.1 Number theory (modular, Chinese remainder, linear Diophantine)

```
// This is a collection of useful code for solving problems that
// involve modular linear equations. Note that all of the
// algorithms described here work on nonnegative integers.
#include <iostream>
#include <vector>
#include <algorithm>
using namespace std;
typedef vector<int> VI;
typedef pair<int, int> PII;
// return a % b (positive value)
```

```
int mod(int a, int b) {
        return ((a%b) + b) % b;
// computes gcd(a,b)
int gcd(int a, int b) {
        while (b) { int t = a%b; a = b; b = t; }
        return a;
// computes lcm(a,b)
int lcm(int a, int b) {
        return a / gcd(a, b) *b;
// (a^b) mod m via successive squaring
int powermod(int a, int b, int m)
        int ret = 1;
        while (b)
                if (b & 1) ret = mod(ret*a, m);
                a = mod(a*a, m);
                b >>= 1:
        return ret;
// returns g = gcd(a, b); finds x, y such that d = ax + by
int extended_euclid(int a, int b, int &x, int &y) {
        int xx = y = 0;
int yy = x = 1;
        while (b) {
                int q = a / b;
                int t = b; b = a%b; a = t;
                t = xx; xx = x - q*xx; x = t;
                t = yy; yy = y - q*yy; y = t;
        return a;
// finds all solutions to ax = b \pmod{n}
VI modular_linear_equation_solver(int a, int b, int n) {
       int x, y;
        VI ret;
        int g = extended_euclid(a, n, x, y);
        if (!(b%g)) {
                   = mod(x*(b / g), n);
                for (int i = 0; i < g; i++)
                        ret.push_back(mod(x + i*(n / g), n));
        return ret:
// computes b such that ab = 1 \pmod{n}, returns -1 on failure
int mod inverse(int a. int n) {
        int x, y;
        int g = extended_euclid(a, n, x, y);
        if (g > 1) return -1;
        return mod(x, n);
// Chinese remainder theorem (special case): find z such that
// z % m1 = r1, z % m2 = r2. Here, z is unique modulo M = lcm(m1, m2).
// Return (z, M). On failure, M = -1
PII chinese_remainder_theorem(int m1, int r1, int m2, int r2) {
        int s, t;
        int g = extended_euclid(m1, m2, s, t);
        if (r1%g != r2%g) return make_pair(0, -1);
        return make_pair(mod(s*r2*m1 + t*r1*m2, m1*m2) / g, m1*m2 / g);
// Chinese remainder theorem: find z such that
// z % m[i] = r[i] for all i. Note that the solution is
// unique modulo M = lcm_i (m[i]). Return (z, M). On
// failure, M = -1. Note that we do not require the a[i]'s
// to be relatively prime.
PII chinese_remainder_theorem(const VI &m, const VI &r) {
        PII ret = make_pair(r[0], m[0]);
        for (int i = 1; i < m.size(); i++) {</pre>
                ret = chinese_remainder_theorem(ret.second, ret.first, m[i], r[i]);
                if (ret.second == -1) break;
        return ret:
// computes x and y such that ax + by = c
// returns whether the solution exists
bool linear_diophantine(int a, int b, int c, int &x, int &y) {
                if (c) return false;
```

```
x = 0; y = 0;
                    return true;
          if (!a)
                    if (c % b) return false;
                    x = 0; y = c / b;
                    return true;
          if (!b)
                    if (c % a) return false;
x = c / a; y = 0;
                    return true:
          int g = gcd(a, b);
          if (c % g) return false;
          x = c / g * mod_inverse(a / g, b / g);
          y = (c - a*x) / b;
          return true;
int main() {
          // expected: 2
          cout << gcd(14, 30) << endl;
          // expected: 2 -2 1
          int x, y;
         int g = extended_euclid(14, 30, x, y);
cout << g << " " << x << " " << y << endl;</pre>
          VI sols = modular_linear_equation_solver(14, 30, 100);
          for (int i = 0; i < sols.size(); i++) cout << sols[i] << " ";</pre>
          cout << endl;</pre>
          // expected: 8
          cout << mod_inverse(8, 9) << endl;</pre>
          // expected: 23 105
         /// II 12
PII ret = chinese_remainder_theorem(VI({ 3, 5, 7 }), VI({ 2, 3, 2 }));
cout < ret.first << " " << ret.second << endl;
         ret = chinese_remainder_theorem(VI({ 4, 6 }), VI({ 3, 5 }));
cout << ret.first << " " << ret.second << endl;
         if (!linear_diophantine(7, 2, 5, x, y)) cout << "ERROR" << endl;
cout << x << " " << y << endl;</pre>
          return 0;
```

3.2 Systems of linear equations, matrix inverse, determinant

```
// Gauss-Jordan elimination with full pivoting.
    (1) solving systems of linear equations (AX=B)
    (2) inverting matrices (AX=I)
    (3) computing determinants of square matrices
// Running time: O(n^3)
            a[][] = an nxn matrix
            b[][] = an nxm matrix
// OUTPUT: X
                   = an nxm matrix (stored in b[][])
            A^{-1} = an nxn matrix (stored in a[][])
            returns determinant of a[][]
#include <iostream>
#include <vector>
#include <cmath>
using namespace std;
const double EPS = 1e-10;
typedef vector<int> VI;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
T GaussJordan(VVT &a, VVT &b) {
  const int n = a.size();
```

```
const int m = b[0].size();
   VI irow(n), icol(n), ipiv(n);
   for (int i = 0; i < n; i++) {</pre>
    int pj = -1, pk = -1;
for (int j = 0; j < n; j++) if (!ipiv[j])</pre>
       for (int k = 0; k < n; k++) if (!ipiv[k])</pre>
    if (pj == -1 || fabs(a[j][k]) > fabs(a[pj][pk])) { pj = j; pk = k; }
if (fabs(a[pj][pk]) < EPS) { cerr << "Matrix is singular." << endl; exit(0); }</pre>
     ipiv[pk]++;
     swap(a[pj], a[pk]);
     swap(b[pj], b[pk]);
    if (pj != pk) det *= -1;
irow[i] = pj;
     icol[i] = pk;
     T c = 1.0 / a[pk][pk];
     det *= a[pk][pk];
     a[pk][pk] = 1.0;
     for (int p = 0; p < n; p++) a[pk][p] *= c;
     for (int p = 0; p < m; p++) b[pk][p] *= c;
     for (int p = 0; p < n; p++) if (p != pk) {
       c = a[p][pk];
       a[p][pk] = 0;
       for (int q = 0; q < n; q++) a[p][q] -= a[pk][q] * c;
for (int q = 0; q < m; q++) b[p][q] -= b[pk][q] * c;
  for (int p = n-1; p >= 0; p--) if (irow[p] != icol[p]) {
    for (int k = 0; k < n; k++) swap(a[k][irow[p]], a[k][icol[p]]);</pre>
  return det;
int main() {
  const int n = 4:
  const int m = 2;
  double A[n][n] = { {1,2,3,4},{1,0,1,0},{5,3,2,4},{6,1,4,6} };
double B[n][m] = { {1,2},{4,3},{5,6},{8,7} };
  VVT a(n), b(n);
  for (int i = 0; i < n; i++) {
    a[i] = VT(A[i], A[i] + n);
    b[i] = VT(B[i], B[i] + m);
  double det = GaussJordan(a, b);
  // expected: 60
  cout << "Determinant: " << det << endl;</pre>
   // expected: -0.233333 0.166667 0.133333 0.0666667
                  0.166667 0.166667 0.333333 -0.333333
                  0.233333 0.833333 -0.133333 -0.0666667
                  0.05 -0.75 -0.1 0.2
  cout << "Inverse: " << endl;</pre>
  for (int i = 0; i < n; i++) {
    for (int j = 0; j < n; j++)
cout << a[i][j] << ' ';
     cout << endl;
   // expected: 1.63333 1.3
                  -0.166667 0.5
                  2.36667 1.7
                 -1.85 -1.35
  cout << "Solution: " << endl;</pre>
  for (int i = 0; i < n; i++) {
  for (int j = 0; j < m; j++) {
    cout << b[i][j] << ' ';
     cout << endl;
```

3.3 Reduced row echelon form, matrix rank

```
// Reduced row echelon form via Gauss-Jordan elimination
// with partial pivoting. This can be used for computing
// the rank of a matrix.
//
// Running time: O(n^3)
//
// INPUT: a[][] = an nxm matrix
//
// OUTPUT: rref[][] = an nxm matrix (stored in a[][])
```

```
returns rank of a[][]
#include <iostream>
#include <vector>
#include <cmath>
using namespace std;
const double EPSILON = 1e-10;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
int rref(VVT &a) {
  int n = a.size();
  int m = a[0].size();
  int r = 0;
  for (int c = 0; c < m && r < n; c++) {
    int j = r;
    for (int i = r + 1; i < n; i++)
     if (fabs(a[i][c]) > fabs(a[j][c])) j = i;
    if (fabs(a[j][c]) < EPSILON) continue;</pre>
    swap(a[j], a[r]);
    T s = 1.0 / a[r][c];
    for (int j = 0; j < m; j++) a[r][j] *= s;
for (int i = 0; i < n; i++) if (i != r) {</pre>
      T t = a[i][c];
      for (int j = 0; j < m; j++) a[i][j] -= t * a[r][j];
    r++;
  return r;
int main() {
  const int n = 5, m = 4;
  double A[n][m] = {
    {16, 2, 3, 13},
    { 5, 11, 10, 8},
    { 9, 7, 6, 12}, { 4, 14, 15, 1}, {13, 21, 21, 13}};
  VVT a(n);
  for (int i = 0; i < n; i++)
    a[i] = VT(A[i], A[i] + m);
  int rank = rref(a);
  // expected: 3
  cout << "Rank: " << rank << endl;
  // expected: 1 0 0 1
               0 1 0 3
                0 0 1 -3
                0 0 0 3.10862e-15
                0 0 0 2.22045e-15
  cout << "rref: " << endl;
  for (int i = 0; i < 5; i++)
    for (int j = 0; j < 4; j++)
cout << a[i][j] << '';
    cout << endl;
```

3.4 Fast Fourier transform

```
#include <cassert>
#include <cstdio>
#include <cstdio>
#include <cmath>

struct cpx
{
    cpx() {}
    cpx(double aa):a(aa),b(0) {}
    cpx(double aa, double bb):a(aa),b(bb) {}
    double a;
    double b;
    double modsq(void) const
    {
        return a * a + b * b;
    }
    cpx bar(void) const
    {
        return cpx(a, -b);
    }
}
```

```
cpx operator + (cpx a, cpx b)
    return cpx(a.a + b.a, a.b + b.b);
cpx operator *(cpx a, cpx b)
    return cpx(a.a * b.a - a.b * b.b, a.a * b.b + a.b * b.a);
cpx operator / (cpx a, cpx b)
    cpx r = a * b.bar():
    return cpx(r.a / b.modsq(), r.b / b.modsq());
cpx EXP (double theta)
    return cpx(cos(theta), sin(theta));
const double two_pi = 4 * acos(0);
// in.
                      input array
// out:
                    output array
// step: {SET TO 1} (used internally)
// size: length of the input/output {MUST BE A POWER OF 2}
// dir: either plus or minus one (direction of the FFT)
// RESULT: \operatorname{out}[k] = \operatorname{\sum_{j=0}^{k}} (j=0) \cdot (size - 1) \cdot \inf[j] * \exp(\operatorname{dir} * 2pi * i * j * k / size)
void FFT(cpx *in, cpx *out, int step, int size, int dir)
    if(size < 1) return;</pre>
    if(size == 1)
        out[0] = in[0];
        return;
    FFT(in, out, step * 2, size / 2, dir);
FFT(in + step, out + size / 2, step * 2, size / 2, dir);
for(int i = 0; i < size / 2; i++)</pre>
        cpx even = out[i];
        cpx odd = out[i + size / 2];
out[i] = even + EXP(dir * two pi * i / size) * odd;
        out[i + size / 2] = even + EXP(dir * two_pi * (i + size / 2) / size) * odd;
// Usage:
// f[0...N-1] and g[0...N-1] are numbers
// Want to compute the convolution h, defined by
// Wall to compute the convocation n, which is n = 1 and n = 1. If n = 1 and n =
// The convolution theorem says H[n] = F[n]G[n] (element-wise product).
// To compute h[] in O(N log N) time, do the following:

    Compute F and G (pass dir = 1 as the argument).

         2. Get H by element-wise multiplying F and G.
         3. Get h by taking the inverse FFT (use dir = -1 as the argument)
                and *dividing by N*. DO NOT FORGET THIS SCALING FACTOR.
int main (void)
    printf("If rows come in identical pairs, then everything works.\n");
    cpx \ a[8] = \{0, 1, cpx(1,3), cpx(0,5), 1, 0, 2, 0\};
    cpx b[8] = \{1, cpx(0,-2), cpx(0,1), 3, -1, -3, 1, -2\};
    cpx A[8];
    cpx B[8];
    FFT(a, A, 1, 8, 1);
    FFT(b, B, 1, 8, 1);
    for(int i = 0; i < 8; i++)
        printf("%7.21f%7.21f", A[i].a, A[i].b);
     printf("\n");
    for (int i = 0; i < 8; i++)
         cpx Ai(0,0);
        for (int j = 0; j < 8; j++)
            Ai = Ai + a[j] \star EXP(j \star i \star two_pi / 8);
        printf("%7.21f%7.21f", Ai.a, Ai.b);
    printf("\n");
     cpx AB[8];
    for(int i = 0; i < 8; i++)
        AB[i] = A[i] * B[i];
```

```
cpx aconvb[8];
FFT(AB, aconvb, 1, 8, -1);
for(int i = 0; i < 8; i++)
    aconvb[i] = aconvb[i] / 8;
for(int i = 0; i < 8; i++)
{
    printf("%7.21f%7.21f", aconvb[i].a, aconvb[i].b);
}
printf("\n");
for(int i = 0; i < 8; i++)
{
    cpx aconvbi(0,0);
    for(int j = 0; j < 8; j++)
    {
        aconvbi = aconvbi + a[j] * b[(8 + i - j) % 8];
    }
    printf("%7.21f%7.21f", aconvbi.a, aconvbi.b);
}
printf("\n");
return 0;</pre>
```

3.5 Simplex algorithm

```
// Two-phase simplex algorithm for solving linear programs of the form
               maximize
               subject to Ax <= b
 // INPUT: A -- an m x n matrix
                  b -- an m-dimensional vector
                      c -- an n-dimensional vector
                      x \ \text{--}\ a vector where the optimal solution will be stored
 // OUTPUT: value of the optimal solution (infinity if unbounded
                        above, nan if infeasible)
// To use this code, create an LPSolver object with A, b, and c as
// arguments. Then, call Solve(x).
#include <iostream>
#include <iomanip>
 #include <vector>
 #include <cmath>
#include <limits>
using namespace std;
typedef long double DOUBLE;
typedef vector<DOUBLE> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
const DOUBLE EPS = 1e-9;
struct LPSolver {
    int m, n;
     VI B, N;
     VVD D;
     LPSolver(const VVD &A, const VD &b, const VD &c) :
         m(b.size()), n(c.size()), N(n + 1), B(m), D(m + 2, VD(n + 2)) {
        for (int i=0, i< m, i++) for (int j=0, j< n, j++) D[i][j]=A[i][j]; for (int i=0, i< m, i++) B[i]=n+i, D[i][n]=-1, D[i][n+1]=b[i], f for (int j=0, j< n, j++) A[i]=j, A[i]=i, A[i
        N[n] = -1; D[m + 1][n] = 1;
     void Pivot(int r, int s) {
         double inv = 1.0 / D[r][s];
         for (int i = 0; i < m + 2; i++) if (i != r)
           for (int j = 0; j < n + 2; j++) if (j != s)
                D[i][j] -= D[r][j] * D[i][s] * inv;
        for (int j = 0; j < n + 2; j++) if (j != s) D[r][j] *= inv;
for (int i = 0; i < m + 2; i++) if (i != r) D[i][s] *= -inv;</pre>
         D[r][s] = inv;
         swap(B[r], N[s]);
    bool Simplex(int phase) {
  int x = phase == 1 ? m + 1 : m;
         while (true) {
            int s = -1;
             for (int j = 0; j \le n; j++) {
                 if (phase == 2 && N[j] == -1) continue;
```

```
if (s == -1 \mid | D[x][j] < D[x][s] \mid | D[x][j] == D[x][s] && N[j] < N[s]) &= j;
      if (D[x][s] > -EPS) return true;
      int r = -1;
      for (int i = 0; i < m; i++) {</pre>
       if (D[i][s] < EPS) continue;</pre>
        if (r == -1 || D[i][n + 1] / D[i][s] < D[r][n + 1] / D[r][s] ||
          if (r == -1) return false;
      Pivot(r, s);
 DOUBLE Solve(VD &x) {
  int r = 0;
    for (int i = 1; i < m; i++) if (D[i][n+1] < D[r][n+1]) r = i;
    if (D[r][n + 1] < -EPS) {
      if (!Simplex(1) || D[m + 1][n + 1] < -EPS) return -numeric_limits<DOUBLE>::infinity();
      for (int i = 0; i < m; i++) if (B[i] == -1) {
        int s = -1;
        for (int j = 0; j \le n; j++)
         if (s == -1 || D[i][j] < D[i][s] || D[i][j] == D[i][s] && N[j] < N[s]) s = j;
    if (!Simplex(2)) return numeric_limits<DOUBLE>::infinity();
    x = VD(n);
    for (int i = 0; i < m; i++) if (B[i] < n) x[B[i]] = D[i][n + 1];</pre>
    return D[m][n + 1];
};
int main() {
  const int m = 4;
  const int n = 3;
  DOUBLE A[m][n] = +
   { 6, -1, 0 },
    \{ -1, -5, 0 \},
   { 1, 5, 1 },
    \{-1, -5, -1\}
  DOUBLE b[m] = \{ 10, -4, 5, -5 \};
  DOUBLE _{c[n]} = \{ 1, -1, 0 \};
  VVD A(m);
  VD b(\underline{b}, \underline{b} + m);
  VD c(\underline{c}, \underline{c} + n);
  for (int i = 0; i < m; i++) A[i] = VD(\_A[i], \_A[i] + n);
  LPSolver solver(A, b, c);
 DOUBLE value = solver.Solve(x):
  cerr << "VALUE: " << value << endl; // VALUE: 1.29032
  cerr << "SOLUTION:"; // SOLUTION: 1.74194 0.451613 1
  for (size_t i = 0; i < x.size(); i++) cerr << " " << x[i];</pre>
  cerr << endl;
 return 0;
```

4 Graph algorithms

4.1 Floyd's algorithm (C++)

```
#include <bits/stdc++.h>
using namespace std;

typedef double T;
typedef vector<T> VT;
typedef vector<VVVT;

typedef vector<VI> VVI;

typedef vector<VI> VVI;

// This function runs the Floyd-Warshall algorithm for all-pairs
// shortest paths. Also handles negative edge weights. Returns true
// if a negative weight cycle is found.
/// Running time: O(|V|^3)
///
```

4.2 Fast Dijkstra's algorithm

```
// Implementation of Dijkstra's algorithm using adjacency lists
// and priority queue for efficiency.
 // Running time: O(|E| log |V|)
 #include <queue>
 #include <cstdio>
using namespace std;
const int INF = 2000000000;
typedef pair<int, int> PII;
int main() {
                                     int N. s. t:
                                     scanf("%d%d%d", &N, &s, &t);
                                      vector<vector<PII> > edges(N);
                                      for (int i = 0; i < N; i++) {
                                                                            scanf("%d", &M);
                                                                           for (int j = 0; j < M; j++) {
                                                                                                                int vertex, dist;
                                                                                                                 scanf("%d%d", &vertex, &dist);
                                                                                                                  edges[i].push_back(make_pair(dist, vertex)); // note order of arguments here
                                      // use priority queue in which top element has the "smallest" priority % \left( 1\right) =\left( 1\right) \left( 
                                    priority_queue<PII, vector<PII>, greater<PII> > Q;
vector(int) dist(N, INF), dad(N, -1);
                                     Q.push(make_pair(0, s));
                                        dist[s] = 0;
                                      while (!Q.empty()) {
                                                                          PII p = Q.top();
                                                                            Q.pop();
                                                                           int here = p.second;
                                                                           if (here == t) break;
                                                                           if (dist[here] != p.first) continue;
                                                                           for (vector<PII>::iterator it = edges[here].begin(); it != edges[here].end(); it++) {
                                                                                                                  if (dist[here] + it->first < dist[it->second]) {
                                                                                                                                                       dist[it->second] = dist[here] + it->first;
                                                                                                                                                        dad[it->second] = here;
                                                                                                                                                        Q.push(make_pair(dist[it->second], it->second));
                                        printf("%d\n", dist[t]);
                                      if (dist[t] < INF)</pre>
                                                                           for (int i = t; i != -1; i = dad[i])
                                                                                                                 printf("%d%c", i, (i == s ? '\n' : ' '));
                                     return 0:
Sample input: 5 0 4
```

```
3 1 4 3 3 4 1
2 0 1 2 3
2 1 5 2 1

Expected:
5 4 2 3 0
*/
```

4.3 Dijkstra's algorithm

```
#include <bits/stdc++.h>
using namespace std;
// (u)int64, float
template<typename D>
struct Dijkstra {
    static constexpr D Inf = 11 << 60;</pre>
    struct Edge { size_t to; D len; };
    struct Vertex {
        vector<Edge> outs;
        D dist = Inf:
#ifdef DIJKSTRA_RECORD_ROUTE
        size_t prev = -1;
#endif
    };
    size_t N;
    vector<Vertex> vs;
    // n nodes
    Dijkstra(size_t n) : N(n) {
        vs.resize(n);
    void add(size_t from, size_t to, D len) {
        assert (from < N):
        assert (to < N);
        assert(len >= 0);
        vs[from].outs.push_back({ to, len });
    void add_u(size_t a, size_t b, D len) {
        add(a, b, len);
        add(b, a, len);
    D solve(size_t from, size_t to) {
        vs[from].dist = 0:
#ifdef DIJKSTRA_RECORD_ROUTE
        vs[from].prev = from;
#endif
        auto comp = [&](size_t x, size_t y) {
    return vs[x].dist < vs[y].dist || (vs[x].dist == vs[y].dist && x < y);</pre>
        set<size_t, decltype(comp) > q { comp };
for (size_t i = 0; i < N; ++i) {</pre>
            q.insert(i);
        while (!q.empty()) {
             size_t i;
                 auto it = q.begin();
                        = *it;
                 q.erase(it);
             if (i == to) {
                 goto RETURN;
             for (Edge const& e : vs[i].outs) {
                 if (vs[e.to].dist > vs[i].dist + e.len) {
                     if (q.find(e.to) != q.end()) {
                          q.erase(e.to);
                          vs[e.to].dist = vs[i].dist + e.len;
#ifdef DIJKSTRA_RECORD_ROUTE
                          vs[e.to].prev = i;
#endif
                          q.insert(e.to);
```

```
RETURN:
    return vs[to].dist;
};
```

#include <bits/stdc++.h>

4.4 Minimum spanning trees using Prim

```
using namespace std;
// (u)int64, float
template<typename D>
struct Prim {
    static constexpr D Inf = 11 << 60;
    struct Edge { size_t to; D len; };
    struct Vertex {
        vector<Edge> outs;
        D dist = Inf;
        size_t prev = -1;
    size_t N;
    vector<Vertex> vs;
    // n nodes
    Prim(size_t n) : N(n) {
        vs.resize(n);
    void add(size_t from, size_t to, D len) {
        assert (from < N);</pre>
        assert (to < N);
        assert (len >= 0);
        vs[from].outs.push_back({ to, len });
    void add_u(size_t a, size_t b, D len) {
        add(a, b, len);
        add(b, a, len);
    D solve(size_t from) {
        vs[from].dist = 0;
        vs[from].prev = from;
        auto comp = [&](size_t x, size_t y) {
            return vs[x].dist < vs[y].dist || (vs[x].dist == vs[y].dist && x < y);</pre>
        set<size_t, decltype(comp)> q { comp };
for (size_t i = 0; i < N; ++i) {</pre>
            q.insert(i);
        while (!q.empty()) {
            size_t i;
                 auto it = q.begin();
                         = *it;
                 q.erase(it);
             if (vs[i].dist == Inf) {
                 vs[i].dist = 0;
             for (Edge const& e : vs[i].outs) {
                 if (vs[e.to].dist > e.len) {
                     if (q.find(e.to) != q.end()) {
                          q.erase(e.to);
                          vs[e.to].dist = e.len;
                          vs[e.to].prev = i;
                          q.insert(e.to);
        D result = 0;
for (size_t i = 0; i < N; ++i)
    result += vs[i].dist;</pre>
        return result;
};
```

4.5 Minimum spanning trees using Kruskal

```
#include <algorithm>
#include <cassert>
 #include <cstdint>
#include <set>
#include <unordered_map>
#include <vector>
using namespace std;
 // (u)int64, float
template <typename D>
struct Kruskal {
          static constexpr D Inf = 11 << 60;</pre>
          struct Edge {
                     size_t n1, n2;
                    D len:
                    bool operator<(Edge const& o) const noexcept {</pre>
                               return len < o.len || (len == o.len && (n1 < o.n1 || (n1 == o.n1 && (n2 < o.n2))));
          };
          typedef unordered_map<size_t, vector<Edge>> Forest;
          vector<Edge> edges;
           // n nodes
          Kruskal(size_t n) : N(n) {}
          void add(size_t n1, size_t n2, D len) {
                     assert (n1 < N);
                     assert (n2 < N);
                     assert(len >= 0);
                     if (n1 > n2) swap(n1, n2);
                     edges.push_back({ n1, n2, len });
          void solve(Forest& forest) {
                     sort(edges.begin(), edges.end());
                     vector<size_t> _ufs(N);
                     for (size_t i = 0; i < N; i++)
    _ufs[i] = i;</pre>
                     auto ufs_p = [&](size_t i) -> size_t* {
                                auto p = i;
                                while (_ufs[p] != p)
                                 p = _ufs[p];
_ufs[i] = p;
                                return &_ufs[i];
                     for (Edge const& e : edges) {
                               auto* ufs_n1p = ufs_p(e.n1);
auto* ufs_n2p = ufs_p(e.n2);
                                if (*ufs_n1p != *ufs_n2p) {
                                         if (*ufs_n1p > *ufs_n2p)
                                                    swap(*ufs_n1p, *ufs_n2p);
                                          auto it = forest.find(*ufs_n2p);
                                          if (it != forest.end()) {
                                                     forest[*ufs\_nlp].insert(forest[*ufs\_nlp].end(), it->second.begin(), it->second.end(), it->second.end
                                                                     ());
                                                     forest.erase(it);
                                          forest[*ufs_n1p].push_back(e);
                                          *ufs_n2p = *ufs_n1p;
                     return:
1:
```

4.6 Strongly connected components

```
#include<memory.h>
struct edge{int e, nxt;};
```

```
int V, E;
edge e[MAXE], er[MAXE];
int sp[MAXV], spr[MAXV];
int group_cnt, group_num[MAXV];
bool v[MAXV];
int stk[MAXV];
void fill_forward(int x)
  int i;
  v[x]=true;
  for(i=sp[x];i;i=e[i].nxt) if(!v[e[i].e]) fill_forward(e[i].e);
  stk[++stk[0]]=x;
void fill backward(int x)
  int i;
 v[x]=false;
  group_num[x]=group_cnt;
  for(i=spr[x];i;i=er[i].nxt) if(v[er[i].e]) fill_backward(er[i].e);
void add_edge(int v1, int v2) //add edge v1->v2
  e [++E].e=v2; e [E].nxt=sp [v1]; sp [v1]=E;
  er[ E].e=v1; er[E].nxt=spr[v2]; spr[v2]=E;
void SCC()
  int i:
  stk[0]=0;
  memset (v. false, sizeof(v));
  for(i=1;i<=V;i++) if(!v[i]) fill_forward(i);</pre>
  for(i=stk[0];i>=1;i--) if(v[stk[i]]){group_cnt++; fill_backward(stk[i]);}
```

4.7 Eulerian path

```
typedef list<Edge>::iterator iter;
struct Edge
        int next_vertex;
        iter reverse_edge;
        Edge(int next_vertex)
               :next_vertex(next_vertex)
};
const int max_vertices = ;
list<Edge> adj[max_vertices];
                                        // adjacency list
vector<int> path;
void find_path(int v)
        while(adj[v].size() > 0)
                int vn = adj[v].front().next_vertex;
                adj[vn].erase(adj[v].front().reverse_edge);
                adj[v].pop_front();
                find_path(vn);
        path.push_back(v);
void add_edge(int a, int b)
        adj[a].push_front(Edge(b));
        iter ita = adj[a].begin();
        adj[b].push_front(Edge(a));
        iter itb = adj[b].begin();
        ita->reverse edge = itb:
        itb->reverse_edge = ita;
```

5 Data structures

5.1 Suffix array

```
// Suffix array construction in O(L log^2 L) time. Routine for
// computing the length of the longest common prefix of any two
// suffixes in O(log L) time.
// TNPUT: string s
// OUTPUT: array suffix[] such that suffix[i] = index (from 0 to L-1)
            of substring s[i...L-1] in the list of sorted suffixes.
            That is, if we take the inverse of the permutation suffix[],
            we get the actual suffix array.
#include <vector>
#include <iostream>
#include <string>
using namespace std;
struct SuffixArray {
  const int L;
  string s;
  vector<vector<int> > P:
  vector<pair<pair<int,int>,int> > M;
  SuffixArray(const string &s) : L(s.length()), s(s), P(1, vector<int>(L, 0)), M(L) {
    for (int i = 0; i < L; i++) P[0][i] = int(s[i]);
for (int skip = 1, level = 1; skip < L; skip *= 2, level++) {
      P.push_back(vector<int>(L, 0));
      for (int i = 0; i < L; i++)
       M[i] = make_pair(make_pair(P[level-1][i], i + skip < L ? P[level-1][i + skip] : -1000), i);
      sort(M.begin(), M.end());
      for (int i = 0; i < L; i++)
         P[level][M[i].second] = (i > 0 \&\& M[i].first == M[i-1].first) ? P[level][M[i-1].second] : i; 
  vector<int> GetSuffixArray() { return P.back(); }
  // returns the length of the longest common prefix of s[i...L-1] and s[j...L-1]
  int LongestCommonPrefix(int i, int j) {
    int len = 0;
    if (i == j) return L - i;
    for (int k = P.size() - 1; k >= 0 && i < L && j < L; k--) {
      if (P[k][i] == P[k][j]) {
        i += 1 << k;
        j += 1 << k;
        len += 1 << k;
    return len;
};
// BEGIN CUT
// The following code solves UVA problem 11512: GATTACA.
#define TESTING
#ifdef TESTING
int main() {
 int T;
  for (int caseno = 0; caseno < T; caseno++) {</pre>
    string s;
    cin >> s:
    SuffixArray array(s);
vector<int> v = array.GetSuffixArray();
    int bestlen = -1, bestpos = -1, bestcount = 0;
    for (int i = 0; i < s.length(); i++) {</pre>
     int len = 0, count = 0;
      for (int j = i+1; j < s.length(); j++) {</pre>
        int 1 = array.LongestCommonPrefix(i, j);
        if (1 >= len) {
          if (1 > len) count = 2; else count++;
          len = 1;
      if (len > bestlen || len == bestlen && s.substr(bestpos, bestlen) > s.substr(i, len)) {
        bestlen = len;
        bestcount = count:
        bestpos = i;
    if (bestlen == 0) {
     cout << "No repetitions found!" << endl;
```

```
} else {
     cout << s.substr(bestpos, bestlen) << " " << bestcount << endl;</pre>
#else
// END CUT
int main() {
  // bobocel is the O'th suffix
  // obocel is the 5'th suffix
      bocel is the 1'st suffix
       ocel is the 6'th suffix
        cel is the 2'nd suffix
         el is the 3'rd suffix
          l is the 4'th suffix
  SuffixArray suffix("bobocel");
  vector<int> v = suffix.GetSuffixArray();
  // Expected output: 0 5 1 6 2 3 4
  for (int i = 0; i < v.size(); i++) cout << v[i] << " ";
  cout << endl:
  cout << suffix.LongestCommonPrefix(0, 2) << endl;</pre>
// BEGIN CUT
#endif
// END CUT
```

5.2 Binary Indexed Tree

```
#include <iostream>
using namespace std;
#define LOGSZ 17
int tree[(1<<LOGSZ)+1];</pre>
int N = (1 << LOGSZ);
// add v to value at x
void set(int x, int v) {
  while (x <= N) {
    tree[x] += v;
    x += (x & -x);
// get cumulative sum up to and including x
int get(int x) {
  int res = 0;
  while(x) {
    res += tree[x]:
    x -= (x & -x);
// get largest value with cumulative sum less than or equal to x;
// for smallest, pass x-1 and add 1 to result
int getind(int x) {
  int idx = 0, mask = N;
  while (mask && idx < N) {
    int t = idx + mask;
    if(x >= tree[t]) {
     idx = t:
     x -= tree[t];
    mask >>= 1;
  return idx;
```

5.3 Union-find set

```
#include <bits/stdc++.h>
using namespace std;
struct UnionFind {
    vector(int> C;
    UnionFind(int n) : C(n) { for (int i = 0; i < n; i++) C[i] = i; }
    int find(int x) { return (C[x] == x) ? x : C[x] = find(C[x]); }
    void merge(int x, int y) { C[find(x)] = find(y); }
}</pre>
```

5.4 KD-tree

```
// A straightforward, but probably sub-optimal KD-tree implmentation
// that's probably good enough for most things (current it's a
// 2D-tree)
// - constructs from n points in O(n \ 1g^2 \ n) time
// - handles nearest-neighbor query in O(lg n) if points are well
    distributed
// - worst case for nearest-neighbor may be linear in pathological
// Sonny Chan, Stanford University, April 2009
#include <iostream>
#include <vector>
#include <limits>
#include <cstdlib>
using namespace std;
// number type for coordinates, and its maximum value
typedef long long ntype;
const ntype sentry = numeric_limits<ntype>::max();
// point structure for 2D-tree, can be extended to 3D
struct point {
    ntype x, y;
    point (ntype xx = 0, ntype yy = 0) : x(xx), y(yy) {}
};
bool operator==(const point &a, const point &b)
    return a.x == b.x && a.y == b.y;
// sorts points on x-coordinate
bool on_x(const point &a, const point &b)
    return a.x < b.x:
// sorts points on v-coordinate
bool on_y(const point &a, const point &b)
    return a.y < b.y;
// squared distance between points
ntype pdist2(const point &a, const point &b)
    ntype dx = a.x-b.x, dy = a.y-b.y;
    return dx*dx + dy*dy;
// bounding box for a set of points
struct bbox
    ntype x0, x1, y0, y1;
    bbox() : x0(sentry), x1(-sentry), y0(sentry), y1(-sentry) {}
    // computes bounding box from a bunch of points
    void compute(const vector<point> &v) {
       y0 = min(y0, v[i].y); y1 = max(y1, v[i].y);
    // squared distance between a point and this bbox, 0 if inside
    ntype distance(const point &p) {
       if (p.x < x0) {
           if (p.y < y0)
                               return pdist2(point(x0, y0), p);
           else if (p.y > y1) return pdist2(point(x0, y1), p);
                               return pdist2(point(x0, p.y), p);
        else if (p.x > x1) {
           if (p.y < y0)
                               return pdist2(point(x1, y0), p);
           else if (p.y > y1) return pdist2(point(x1, y1), p);
           else
                               return pdist2(point(x1, p.y), p);
       else
                               return pdist2(point(p.x, y0), p);
           if (p.v < v0)
           else if (p.y > y1) return pdist2(point(p.x, y1), p);
           else
                               return 0;
```

```
};
// stores a single node of the kd-tree, either internal or leaf
struct kdnode
    bool leaf;
                      // true if this is a leaf node (has one point)
    point pt;
                      \ensuremath{//} the single point of this is a leaf
    bbox bound;
                     // bounding box for set of points in children
    kdnode *first, *second; // two children of this kd-node
    kdnode() : leaf(false), first(0), second(0) {}
    "kdnode() { if (first) delete first; if (second) delete second; }
    // intersect a point with this node (returns squared distance)
    ntype intersect(const point &p) {
        return bound.distance(p);
    // recursively builds a kd-tree from a given cloud of points
    void construct(vector<point> &vp)
         // compute bounding box for points at this node
        bound.compute(vp);
         // if we're down to one point, then we're a leaf node
        if (vp.size() == 1) {
             leaf = true;
             pt = vp[0];
        else {
                split on x if the bbox is wider than high (not best heuristic...)
             if (bound.x1-bound.x0 >= bound.y1-bound.y0)
                 sort(vp.begin(), vp.end(), on_x);
             // otherwise split on y-coordinate
             else
                 sort(vp.begin(), vp.end(), on_y);
             /\!/ divide by taking half the array for each child /\!/ (not best performance if many duplicates in the middle)
            int half = vp.size()/2;
vector<point> vl(vp.begin(), vp.begin() +half;
vector<point> vr(vp.begin() +half, vp.end());
first = new kdnode(); first->construct(vl);
             second = new kdnode(); second->construct(vr);
};
// simple kd-tree class to hold the tree and handle queries
struct kdtree
    kdnode +root:
    // constructs a kd-tree from a points (copied here, as it sorts them)
    kdtree(const vector<point> &vp) {
        vector<point> v(vp.begin(), vp.end());
        root = new kdnode();
        root->construct(v);
    "kdtree() { delete root; }
    // recursive search method returns squared distance to nearest point
    ntype search(kdnode *node, const point &p)
        if (node->leaf) {
             // commented special case tells a point not to find itself
              if (p == node->pt) return sentry;
               else
                 return pdist2(p, node->pt);
        ntype bfirst = node->first->intersect(p);
        ntype bsecond = node->second->intersect(p);
         // choose the side with the closest bounding box to search first
         // (note that the other side is also searched if needed)
        if (bfirst < bsecond) {
             ntype best = search(node->first, p);
             if (bsecond < best)</pre>
                 best = min(best, search(node->second, p));
             return best:
        else {
             ntype best = search(node->second, p);
             if (bfirst < best)</pre>
                 best = min(best, search(node->first, p));
             return best;
```

```
// squared distance to the nearest
ntype nearest(const point &p) {
    return search(root, p);
};

//
// some basic test code here

int main() {
    // generate some random points for a kd-tree
    vector<point> vp;
    for (int i = 0; i < 100000; ++i) {
        vp. push_back(point(rand()%100000, rand()%100000));
    }
    kdtree tree(vp);

    // query some points
    for (int i = 0; i < 10; ++i) {
        point q(rand()%100000, rand()%100000);
        cout << "Closest squared distance to (" << q.x << ", " << q.y << ")"
        << " is " << tree.nearest(q) << endl;
}

return 0;
}</pre>
```

5.5 Splay tree

```
#include <cstdio>
#include <algorithm>
using namespace std;
const int N MAX = 130010:
const int oo = 0x3f3f3f3f3f
struct Node
  Node *ch[2], *pre;
  int val. size:
  bool isTurned;
} nodePool[N_MAX], *null, *root;
Node *allocNode(int val)
  static int freePos = 0;
  Node *x = &nodePool[freePos ++];
  x->val = val, x->isTurned = false;
  x->ch[0] = x->ch[1] = x->pre = null;
  x->size = 1;
 return x:
inline void update (Node *x)
  x->size = x->ch[0]->size + x->ch[1]->size + 1;
inline void makeTurned(Node *x)
  if(x == null)
    return;
  swap(x->ch[0], x->ch[1]);
x->isTurned ^= 1;
inline void pushDown(Node *x)
  if(x->isTurned)
    makeTurned(x->ch[0]);
    makeTurned(x->ch[1]);
    x->isTurned ^= 1;
inline void rotate(Node *x, int c)
  Node *y = x->pre;
  x->pre = y->pre;
  if(y->pre != null)
   y->pre->ch[y == y->pre->ch[1]] = x;
    ->ch[!c] = x->ch[c];
  if(x->ch[c] != null)
    x->ch[c]->pre = y;
```

```
x->ch[c] = y, y->pre = x;
  update(y);
  if(y == root)
   root = x;
void splay(Node *x, Node *p)
  while (x->pre != p)
    if(x->pre->pre == p)
      rotate(x, x == x->pre->ch[0]);
    else
      Node *y = x->pre, *z = y->pre;
      if(y == z->ch[0])
        if(x == y->ch[0])
          rotate(y, 1), rotate(x, 1);
          rotate(x, 0), rotate(x, 1);
      else
        if(x == y->ch[1])
          rotate(y, 0), rotate(x, 0);
        else
          rotate(x, 1), rotate(x, 0);
 update(x);
void select(int k, Node *fa)
  Node *now = root;
  while (1)
   pushDown (now);
   int tmp = now->ch[0]->size + 1;
if(tmp == k)
     break:
    else if(tmp < k)</pre>
      now = now -> ch[1], k -= tmp;
    else
     now = now -> ch[0];
  splay(now, fa);
Node *makeTree(Node *p, int 1, int r)
 if(1 > r)
   return null:
  int mid = (1 + r) / 2:
 Node *x = allocNode(mid);
 x->re = p;
x->ch[0] = makeTree(x, 1, mid - 1);
x->ch[1] = makeTree(x, mid + 1, r);
  update(x);
  return x;
int main()
  int n, m;
 null = allocNode(0);
 null->size = 0;
  root = allocNode(0);
  root->ch[1] = allocNode(oo);
  root->ch[1]->pre = root;
  update(root);
  scanf("%d%d", &n, &m);
  root->ch[1]->ch[0] = makeTree(root->ch[1], 1, n);
  splay(root->ch[1]->ch[0], null);
  while (m --)
   int a, b;
   scanf("%d%d", &a, &b);
    a ++, b ++;
    select(a - 1, null);
    select(b + 1, root);
    makeTurned(root->ch[1]->ch[0]);
  for(int i = 1; i <= n; i ++)
    select(i + 1, null);
    printf("%d ", root->val);
```

5.6 Lowest common ancestor

```
const int max_nodes, log_max_nodes;
int num_nodes, log_num_nodes, root;
vector<int> children[max_nodes];
                                         // children[i] contains the children of node i
int A[max_nodes][log_max_nodes+1];
                                         // A[i][j] is the 2^j-th ancestor of node i, or -1 if that
      ancestor does not exist
int L[max_nodes];
                                         // L[i] is the distance between node i and the root
// floor of the binary logarithm of n
int lb(unsigned int n)
   if(n==0)
       return -1;
    int p = 0;
    if (n >= 1<<16) { n >>= 16; p += 16; }
    if (n >= 1 << 8) { n >>= 8; p += 8; }
    if (n >= 1<< 4) { n >>= 4; p += 4; ]
    if (n >= 1 << 2) { n >>= 2; p += 2; }
   if (n >= 1<< 1) {
    return p;
void DFS(int i, int 1)
    for(int j = 0; j < children[i].size(); j++)</pre>
       DFS(children[i][j], 1+1);
int LCA(int p, int q)
    // ensure node p is at least as deep as node q
    if(L[p] < L[q])
        swap(p, q);
    // "binary search" for the ancestor of node p situated on the same level as {\bf q}
    for(int i = log_num_nodes; i >= 0; i--)
        if(L[p] - (1 << i) >= L[q])
            p = A[p][i];
    if(p == q)
        return p;
    // "binary search" for the LCA
    for(int i = log_num_nodes; i >= 0; i--)
       if (A[p][i] != -1 && A[p][i] != A[q][i])
            p = A[p][i];
            q = A[q][i];
    return A[p][0];
int main(int argc,char* argv[])
    // read num_nodes, the total number of nodes
    log_num_nodes=1b(num_nodes);
    for(int i = 0; i < num_nodes; i++)</pre>
        // read p, the parent of node i or -1 if node i is the root
        A[i][0] = p;
        if (p != -1)
            children[p].push_back(i);
            root = i;
    // precompute A using dynamic programming
    for(int j = 1; j <= log_num_nodes; j++)</pre>
        for(int i = 0; i < num_nodes; i++)</pre>
            if(A[i][j-1] != -1)
                A[i][j] = A[A[i][j-1]][j-1];
            else
                A[i][j] = -1;
    // precompute L
    DFS(root, 0);
```

6 Miscellaneous

6.1 Longest increasing subsequence

```
// Given a list of numbers of length n, this routine extracts a
// longest increasing subsequence.
// Running time: O(n log n)
    INPUT: a vector of integers
   OUTPUT: a vector containing the longest increasing subsequence
#include <iostream>
#include <vector>
#include <algorithm>
using namespace std;
typedef vector<int> VI;
typedef pair<int, int> PII;
typedef vector<PII> VPII;
#define STRICTLY_INCREASNG
VI LongestIncreasingSubsequence(VI v) {
 VPII best;
  VI dad(v.size(), -1);
  for (int i = 0; i < v.size(); i++) {</pre>
#ifdef STRICTLY_INCREASIG
   PII item = make_pair(v[i], 0);
    VPII::iterator it = lower_bound(best.begin(), best.end(), item);
    item.second = i;
#else
   PII item = make_pair(v[i], i);
    VPII::iterator it = upper_bound(best.begin(), best.end(), item);
#endif
   if (it == best.end()) {
      dad[i] = (best.size() == 0 ? -1 : best.back().second);
      best.push_back(item);
   } else {
     dad[i] = it == best.begin() ? -1 : prev(it)->second;
      *it = item;
  for (int i = best.back().second; i >= 0; i = dad[i])
   ret.push_back(v[i]);
  reverse(ret.begin(), ret.end());
  return ret:
```

6.2 Dates

```
// Routines for performing computations on dates. In these routines,
// months are expressed as integers from 1 to 12, days are expressed
// as integers from 1 to 31, and years are expressed as 4-digit
// integers.

#include <iostream>
#include <string>
using namespace std;

string dayOfWeek[] = {"Mon", "Tue", "Wed", "Thu", "Fri", "Sat", "Sun"};

// converts Gregorian date to integer (Julian day number)
int dateToInt (int m, int d, int y){
    return
    1461 * (y + 4800 + (m - 14) / 12) / 4 +
    367 * (m - 2 - (m - 14) / 12 + 12) / 12 -
    3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +
    d - 32075;
}

// converts integer (Julian day number) to Gregorian date: month/day/year
```

```
void intToDate (int jd, int &m, int &d, int &y) {
  int x, n, i, j;
  x = jd + 68569;
 n = 4 * x / 146097;
  x = (146097 * n + 3) / 4;
  i = (4000 * (x + 1)) / 1461001;
  x -= 1461 * i / 4 - 31;
  j = 80 * x / 2447;
  d = x - 2447 * j / 80;
 x = j / 11;

m = j + 2 - 12 * x;
 y = 100 * (n - 49) + i + x;
// converts integer (Julian day number) to day of week
string intToDay (int jd) {
 return dayOfWeek[jd % 7];
int main (int argc, char **argv) {
 int jd = dateToInt (3, 24, 2004);
  int m, d, y;
  intToDate (jd, m, d, y);
  string day = intToDay (jd);
  // expected output:
      2453089
       3/24/2004
  // Wed
  cout << jd << endl
   << m << "/" << d << "/" << y << endl
    << day << endl;
```

6.3 Prime numbers

```
// O(sqrt(x)) Exhaustive Primality Test
#include <cmath>
#define EPS 1e-7
typedef long long LL;
bool IsPrimeSlow (LL x)
  if(x<=1) return false;</pre>
  if(x<=3) return true;</pre>
  if (!(x%2) || !(x%3)) return false;
  LL s=(LL) (sqrt((double)(x))+EPS);
  for (LL i=5; i<=s; i+=6)
    if (!(x%i) || !(x%(i+2))) return false;
  return true;
// Primes less than 1000:
                                 59
                                       61
                                             67
                                                   71
                                                                79
                                                                       8.3
                         107
                                109
                                            127
                                                  131
                                                               139
                                                                     149
      157
                                179
                                      181
                                            191
                                                  193
                                                         197
                                                               199
            163
                   167
                                                                     211
                                                  263
                                                         269
                         239
                               241
                                      251
            293
                         311
                                313
                                                  337
                         383
                                389
                                      397
                                            401
                                                  409
      439
            443
                   449
                         457
                                461
                                      463
                                            467
                                                  479
                                                         487
                                                               491
      509
                   523
                         541
                                547
                                      557
                                            563
                                                  569
                                                         571
      599
            601
                   607
                         613
                                617
                                      619
                                            631
                                                  641
                                                         643
                                                               647
      661
            673
                         683
                                691
                                      701
                                            709
                                                  719
                                                         727
                                                                      739
                                                                            743
      751
829
                  761
853
                         769
857
                               773
859
                                      787
863
                                                               821
887
                                            797
                                                  809
                                            877
            839
                                                  881
                                                         883
                                                                     907
                                                                            911
                               947
                                      953
                         941
// Other primes:
      The largest prime smaller than 10 is 7.
      The largest prime smaller than 100 is 97.
      The largest prime smaller than 1000 is 997.
      The largest prime smaller than 10000 is 9973.
      The largest prime smaller than 100000 is 99991
      The largest prime smaller than 1000000 is 999983.
      The largest prime smaller than 10000000 is 9999991.
      The largest prime smaller than 100000000 is 99999989.
      The largest prime smaller than 1000000000 is 999999937.
      The largest prime smaller than 10000000000 is 9999999967. The largest prime smaller than 10000000000 is 99999999977.
      The largest prime smaller than 100000000000 is 99999999999.
The largest prime smaller than 1000000000000 is 999999999971.
The largest prime smaller than 100000000000000 is 999999999973.
      The largest prime smaller than 100000000000000 is 999999999999937
      The largest prime smaller than 100000000000000 is 9999999999999997
```

6.4 C++ input/output

```
#include <bits/stdc++.h>
using namespace std;
int main()
    // Ouput a specific number of digits past the decimal point,
    // in this case 5
    cout.setf(ios::fixed); cout << setprecision(5);</pre>
    cout << 100.0/7.0 << endl;
    cout.unsetf(ios::fixed);
    // Output the decimal point and trailing zeros
    cout.setf(ios::showpoint);
    cout << 100.0 << endl:
    cout.unsetf(ios::showpoint);
    // Output a '+' before positive values
    cout.setf(ios::showpos);
cout << 100 << " " << -100 << endl;
    cout.unsetf(ios::showpos);
    // Output numerical values in hexadecimal
    cout << hex << 100 << " " << 1000 << " " << 10000 << dec << endl;
```

6.5 Knuth-Morris-Pratt

```
Finds all occurrences of the pattern string p within the
text string t. Running time is O(n + m), where n and m
are the lengths of p and t, respecitvely.
#include <iostream>
#include <string>
#include <vector>
using namespace std;
typedef vector<int> VI;
void buildPi(string& p, VI& pi)
  pi = VI(p.length());
  int k = -2;
  int k = -2;
for(int i = 0; i < p.length(); i++) {
  while (k >= -1 && p[k+1] != p[i])
  k = (k == -1) ? -2 : pi[k];
    pi[i] = ++k;
int KMP (string& t, string& p)
  VI pi;
  buildPi(p, pi);
  int k = -1;
  for(int i = 0; i < t.length(); i++) {
  while(k >= -1 && p[k+1] != t[i])
      k = (k == -1) ? -2 : pi[k];
    k++;
    if(k == p.length() - 1) {
      // p matches t[i-m+1, ..., i]
      cout << "matched at index " << i-k << ": ";
       cout << t.substr(i-k, p.length()) << endl;</pre>
       k = (k == -1) ? -2 : pi[k];
  return 0:
int main()
  string a = "AABAACAADAABAABA", b = "AABA";
  KMP(a, b); // expected matches at: 0, 9, 12
  return 0:
```

6.6 Latitude/longitude

```
Converts from rectangular coordinates to latitude/longitude and vice
versa. Uses degrees (not radians).
#include <iostream>
#include <cmath>
using namespace std;
struct 11
  double r, lat, lon;
};
struct rect
  double x, y, z;
};
11 convert (rect& P)
 11 Q;
  Q.r = sqrt(P.x*P.x+P.y*P.y+P.z*P.z);
  Q.lat = 180/M_PI*asin(P.z/Q.r);
  Q.lon = 180/M_PI*acos(P.x/sqrt(P.x*P.x+P.y*P.y));
rect convert(11& Q)
 P.x = Q.r*cos(Q.lon*M_PI/180)*cos(Q.lat*M_PI/180);
P.y = Q.r*sin(Q.lon*M_PI/180)*cos(Q.lat*M_PI/180);
P.z = Q.r*sin(Q.lat*M_PI/180);
  return P;
int main()
  A.x = -1.0; A.y = 2.0; A.z = -3.0;
 B = convert(A);
cout << B.r << " " << B.lat << " " << B.lon << endl;</pre>
  A = convert(B);
  cout << A.x << " " << A.y << " " << A.z << endl;
```

6.7 Vim settings

```
set enc=uff-8
set fenc=utf-8
scriptencoding utf-8
set fencs=utf-8,ucs-bom,gb18030,gbk,gb2312,cp936
syntax on
filetype plugin on
filetype indent on
set mouse=a
set nocompatible
set tabstop=4
set shiftwidth=4
set expandtab
set smarttab
set autoindent
set textwidth=1000
```

```
set showmatch

set ruler

set hlsearch
set incsearch
set ignorecase
set smartcase

set number
set relativenumber

set fdm=marker
set scrolloff=10
set showcmd
set backspace=indent,eol,start
set formatoptions=q,r
set nowrap
set foldmethod=indent
set folddevelstart=10
```

```
set wildmenu
set omnifunc=syntaxcomplete#Complete
"the trail config failed in generate latex
"remember to add in the site"
set list listchars=tab:>-,trail:

set t_Co=256
set background=dark
nnoremap;:
inoremap <silent> jj <ESC>
nnoremap <silent> <HOME> ^
inoremap <silent> <HOME> Ow
nnoremap <silent> <HOME> Ow
nnoremap <silent> <F9> :set nowrap<CR>
nnoremap <silent> <F9> :set nowrap<CR>
nnoremap <silent> <F10> :set wrap<CR>
```