Stanford University ICPC Team Notebook (2015-16)

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1 Combinatorial optimization

1.1 Edmonds-Karp max-flow (normal, SPFA min-cost)

```
#include <bits/stdc++.h>
using namespace std;

template<typename D>
struct EdmondsKarp {
    static constexpr D Inf = 11 << 60;</pre>
```

```
struct Edge {
        size_t from, to; D cap;
    size_t N;
    vector<Edge> edges;
    vector<vector<size_t>> outs;
    EdmondsKarp(size_t n) : N(n) {
        outs.resize(n);
    void add(size_t from, size_t to, D cap) {
        assert (from < N);
         assert (to < N);
        if (cap > 0) {
            outs[from].emplace_back(edges.size());
             edges.emplace_back(Edge { from, to, cap });
             outs[to].emplace_back(edges.size());
             edges.emplace_back(Edge { to, from, 0 });
    vector<D> amount:
    vector<size t> route:
    vector<size_t> visiting;
    D solve(size_t from, size_t to) {
        amount.resize(N);
        route.resize(N);
        visiting.resize(N);
        D flow = 0;
        while (true) {
             fill(amount.begin(), amount.end(), 0);
             size_t head = 0, tail = 0;
             amount[from] = Inf;
             visiting[tail++] = from;
             while (head < tail && !amount[to]) {</pre>
                 size_t i = visiting[head++];
                 for (size_t j = 0; j < outs[i].size(); ++j) {</pre>
                     Edge &e = edges[outs[i][j]];
                     if (!amount[e.to] && e.cap) {
                         amount[e.to] = min(amount[i], e.cap);
route[e.to] = outs[i][j];
                         visiting[tail++] = e.to;
            if (!amount[to]) break;
             for (size_t i = to; i != from; i = edges[route[i]].from) {
                 edges[route[i]].cap -= amount[to];
edges[route[i] ^ 1].cap += amount[to];
             flow += amount[to];
        return flow;
};
```

1.2 Dinic max-flow

```
// Adjacency list implementation of Dinic's blocking flow algorithm.
// This is very fast in practice, and only loses to push-relabel flow.
// Running time:
// O(|V|^2 |E|)
// INPUT:
// - graph, constructed using AddEdge()
// - source and sink
//
// OUTPUT:
// - maximum flow value
// - To obtain actual flow values, look at edges with capacity > 0
// (zero capacity edges are residual edges).
```

```
#include<cstdio>
#include<vector>
#include<queue>
using namespace std;
typedef long long LL;
struct Edge {
 int u, v;
  LL cap, flow;
  Edge () {}
 Edge(int u, int v, LL cap): u(u), v(v), cap(cap), flow(0) {}
struct Dinic (
  int N;
  vector<Edge> E;
  vector<vector<int>> g;
  vector<int> d, pt;
  Dinic(int N): N(N), E(0), g(N), d(N), pt(N) {}
  void AddEdge(int u, int v, LL cap) {
    if (u != v) {
      E.emplace_back(u, v, cap);
      g[u].emplace_back(E.size() - 1);
      E.emplace_back(v, u, 0);
      g[v].emplace_back(E.size() - 1);
  bool BFS(int S, int T) {
    queue<int> q({S});
    fill(d.begin(), d.end(), N + 1);
    while(!q.empty()) {
      int u = q.front(); q.pop();
      if (u == T) break;
      for (int k: g[u]) {
        Edge &e = E[k];
        if (e.flow < e.cap && d[e.v] > d[e.u] + 1) {
  d[e.v] = d[e.u] + 1;
          q.emplace(e.v);
    return d[T] != N + 1;
  LL DFS(int u, int T, LL flow = -1) {
    if (u == T || flow == 0) return flow;
    for (int &i = pt[u]; i < g[u].size(); ++i) {</pre>
      Edge &e = E[g[u][i]];
      Edge &oe = E[g[u][i]^1];
      if(d[e.v] == d[e.u] + 1) {
        LL amt = e.cap - e.flow;
if (flow != -1 && amt > flow) amt = flow;
        if (LL pushed = DFS(e.v, T, amt)) {
          e.flow += pushed;
          oe.flow -= pushed;
          return pushed;
    return 0;
  LL MaxFlow(int S, int T) {
    LL total = 0;
    while (BFS(S, T)) {
      fill(pt.begin(), pt.end(), 0);
while (LL flow = DFS(S, T))
        total += flow;
    return total;
// The following code solves SPOJ problem #4110: Fast Maximum Flow (FASTFLOW)
int main()
  scanf("%d%d", &N, &E);
  Dinic dinic(N);
  for(int i = 0; i < E; i++)</pre>
    int u, v;
    LL cap;
    scanf("%d%d%lld", &u, &v, &cap);
    dinic.AddEdge(u - 1, v - 1, cap);
```

```
dinic.AddEdge(v - 1, u - 1, cap);
}
printf("%lld\n", dinic.MaxFlow(0, N - 1));
return 0;
}
// END CUT
```

1.3 Push-relabel max-flow

```
// Adjacency list implementation of FIFO push relabel maximum flow
// with the gap relabeling heuristic. This implementation is // significantly faster than straight Ford-Fulkerson. It solves
// random problems with 10000 vertices and 1000000 edges in a few
// seconds, though it is possible to construct test cases that
// achieve the worst-case.
// Running time:
      0(|V|^3)
// INPUT:
        - graph, constructed using AddEdge()
       - source
       - sink
       - maximum flow value
       - To obtain the actual flow values, look at all edges with
         capacity > 0 (zero capacity edges are residual edges).
#include <cmath>
#include <vector>
#include <iostream>
#include <queue>
using namespace std;
typedef long long LL;
struct Edge {
 int from, to, cap, flow, index;
  Edge(int from, int to, int cap, int flow, int index) :
    from (from), to (to), cap(cap), flow(flow), index(index) {}
struct PushRelabel {
  int N;
  vector<vector<Edge> > G;
  vector<LL> excess;
  vector<int> dist, active, count;
  queue<int> Q;
  PushRelabel(int N) : N(N), G(N), excess(N), dist(N), active(N), count(2*N) {}
  void AddEdge(int from, int to, int cap) {
   G[from].push_back(Edge(from, to, cap, 0, G[to].size()));
   if (from == to) G[from].back().index++;
    G[to].push_back(Edge(to, from, 0, 0, G[from].size() - 1));
  void Enqueue(int v) {
    if (!active[v] && excess[v] > 0) { active[v] = true; Q.push(v); }
  void Push (Edge &e) {
    int amt = int(min(excess[e.from], LL(e.cap - e.flow)));
    if (dist[e.from] <= dist[e.to] || amt == 0) return;</pre>
    e.flow += amt:
    G[e.to][e.index].flow -= amt;
    excess[e.to] += amt;
    excess[e.from] -= amt;
    Enqueue (e.to):
  void Gap(int k) {
    for (int v = 0; v < N; v++) {
      if (dist[v] < k) continue;</pre>
       count[dist[v]]--;
      dist[v] = max(dist[v], N+1);
      count[dist[v]]++;
      Enqueue (v);
  void Relabel(int v) {
    count[dist[v]]--;
    dist[v] = 2*N;
    for (int i = 0; i < G[v].size(); i++)</pre>
```

```
if (G[v][i].cap - G[v][i].flow > 0)
        dist[v] = min(dist[v], dist[G[v][i].to] + 1);
    count[dist[v]]++;
    Enqueue (v);
  void Discharge(int v) {
    for (int i = 0; excess[v] > 0 && i < G[v].size(); i++) Push(G[v][i]);
    if (excess[v] > 0) {
      if (count[dist[v]] == 1)
        Gap(dist[v]);
      else
        Relabel(v);
  LL GetMaxFlow(int s, int t) {
    count[0] = N-1;
    count[N] = 1;
    dist[s] = N;
    active[s] = active[t] = true;
for (int i = 0; i < G[s].size(); i++) {
  excess[s] += G[s][i].cap;</pre>
      Push(G[s][i]);
    while (!Q.empty()) {
      int v = 0.front();
      Q.pop();
      active[v] = false;
      Discharge(v);
    LL totflow = 0;
    for (int i = 0; i < G[s].size(); i++) totflow += G[s][i].flow;</pre>
    return totflow;
};
// BEGIN CUT
// The following code solves SPOJ problem #4110: Fast Maximum Flow (FASTFLOW)
int main() {
  int n, m;
  scanf("%d%d", &n, &m);
  PushRelabel pr(n);
  for (int i = 0; i < m; i++) {
    scanf("%d%d%d", &a, &b, &c);
    if (a == b) continue;
    pr.AddEdge(a-1, b-1, c);
    pr.AddEdge(b-1, a-1, c);
  printf("%Ld\n", pr.GetMaxFlow(0, n-1));
  return 0:
// END CUT
```

1.4 Min-cost matching

```
// Min cost bipartite matching via shortest augmenting paths
^{\prime\prime} // This is an O(n^3) implementation of a shortest augmenting path
// algorithm for finding min cost perfect matchings in dense
// graphs. In practice, it solves 1000x1000 problems in around 1
// second.
    cost[i][j] = cost for pairing left node i with right node j
    Lmate[i] = index of right node that left node i pairs with
    Rmate[j] = index of left node that right node j pairs with
// The values in cost[i][j] may be positive or negative. To perform
// maximization, simply negate the cost[][] matrix.
#include <algorithm>
#include <cstdio>
#include <cmath>
#include <vector>
using namespace std;
typedef vector<double> VD;
typedef vector<VD> VVD;
```

```
typedef vector<int> VI;
double MinCostMatching(const VVD &cost, VI &Lmate, VI &Rmate) {
  int n = int(cost.size());
   // construct dual feasible solution
  VD u(n);
  VD v(n);
  for (int i = 0; i < n; i++) {</pre>
    u[i] = cost[i][0];
     for (int j = 1; j < n; j++) u[i] = min(u[i], cost[i][j]);</pre>
  for (int j = 0; j < n; j++) {
    v[j] = cost[0][j] - u[0];
    for (int i = 1; i < n; i++) v[j] = min(v[j], cost[i][j] - u[i]);</pre>
   // construct primal solution satisfying complementary slackness
   Lmate = VI(n, -1);
  Rmate = VI(n, -1);
  int mated = 0;
for (int i = 0; i < n; i++) {
   for (int j = 0; j < n; j++) {
      if (Rmate[j] != -1) continue;
}</pre>
       if (fabs(cost[i][j] - u[i] - v[j]) < 1e-10) {</pre>
         Lmate[i] = j;
Rmate[j] = i;
         mated++;
         break:
  VD dist(n);
  VI dad(n);
  VI seen(n);
   // repeat until primal solution is feasible
  while (mated < n) {</pre>
     // find an unmatched left node
     int s = 0:
     while (Lmate[s] != -1) s++;
     // initialize Dijkstra
     fill(dad.begin(), dad.end(), -1);
     fill(seen.begin(), seen.end(), 0);
     for (int k = 0; k < n; k++)
       dist[k] = cost[s][k] - u[s] - v[k];
     int j = 0;
     while (true) {
       // find closest
       \frac{1}{1} = -1:
       for (int k = 0; k < n; k++) {
         if (seen[k]) continue;
         if (j == -1 || dist[k] < dist[j]) j = k;</pre>
       seen[j] = 1;
        // termination condition
       if (Rmate[j] == -1) break;
       // relax neighbors
       const int i = Rmate[j];
for (int k = 0; k < n; k++) {
  if (seen[k]) continue;</pre>
         const double new_dist = dist[j] + cost[i][k] - u[i] - v[k];
         if (dist[k] > new_dist) {
  dist[k] = new_dist;
           dad[k] = j;
     // update dual variables
     for (int k = 0; k < n; k++) {
   if (k == j || !seen[k]) continue;</pre>
       const int i = Rmate[k];
       v[k] += dist[k] - dist[j];
       u[i] -= dist[k] - dist[j];
     u[s] += dist[j];
     // augment along path
     while (dad[j] >= 0) {
      const int d = dad[j];
       Rmate[j] = Rmate[d];
       Lmate[Rmate[j]] = j;
       j = d;
```

```
Rmate[j] = s;
Lmate[s] = j;
mated++;
} }
double value = 0;
for (int i = 0; i < n; i++)
value += cost[i][Lmate[i]];
return value;
```

1.5 Max bipartite matching

```
// This code performs maximum bipartite matching.
// Running time: O(|E| |V|) -- often much faster in practice
     INPUT: w[i][j] = edge between row node i and column node j
     OUTPUT: mr[i] = assignment for row node i, -1 if unassigned
              mc[j] = assignment for column node j, -1 if unassigned
              function returns number of matches made
#include <vector>
using namespace std;
typedef vector<int> VI;
typedef vector<VI> VVI;
bool FindMatch(int i, const VVI &w, VI &mr, VI &mc, VI &seen) {
  for (int j = 0; j < w[i].size(); j++) {
  if (w[i][j] && !seen[j]) {</pre>
      seen[j] = true;
      if (mc[j] < 0 \mid \mid FindMatch(mc[j], w, mr, mc, seen)) {
       mr[i] = j;
mc[i] = i;
        return true;
  return false;
int BipartiteMatching(const VVI &w, VI &mr, VI &mc) {
 mr = VI(w.size(), -1);
  mc = VI(w[0].size(), -1);
  for (int i = 0; i < w.size(); i++) {</pre>
    VI seen(w[0].size());
    if (FindMatch(i, w, mr, mc, seen)) ct++;
  return ct;
```

1.6 Global min-cut

```
// Adjacency matrix implementation of Stoer-Wagner min cut algorithm.
// Running time:
// O(|V|^3)
// INPUT:
// - graph, constructed using AddEdge()
//
// OUTPUT:
// - (min cut value, nodes in half of min cut)
#include <cmath>
#include <vector>
#include <iostream>
using namespace std;
typedef vector<Int> VI;
typedef vector<VI> VVI;
const int INF = 10000000000;
pair<int, VI> GetMinCut(VVI &weights) {
   int N = weights.size();
```

```
VI used(N), cut, best_cut;
  int best_weight = -1;
  for (int phase = N-1; phase >= 0; phase--) {
    VI w = weights[0];
    VI added = used;
    int prev, last = 0;
    for (int i = 0; i < phase; i++) {</pre>
      prev = last;
      last = -1;
      for (int j = 1; j < N; j++)
        if (!added[j] && (last == -1 || w[j] > w[last])) last = j;
      if (i == phase-1) {
        for (int j = 0; j < N; j++) weights[prev][j] += weights[last][j];
for (int j = 0; j < N; j++) weights[j][prev] = weights[prev][j];
used[last] = true;</pre>
        cut.push_back(last);
        if (best_weight == -1 || w[last] < best_weight) {</pre>
          best_cut = cut;
          best_weight = w[last];
      } else {
        for (int j = 0; j < N; j++)
          w[j] += weights[last][j];
        added[last] = true;
 return make_pair(best_weight, best_cut);
// The following code solves UVA problem #10989: Bomb, Divide and Conquer
  cin >> N;
  for (int i = 0; i < N; i++) {
   int n, m;
    cin >> n >> m;
    VVI weights(n, VI(n));
    for (int j = 0; j < m; j++) {
     int a, b, c;
     cin >> a >> b >> c;
      weights[a-1][b-1] = weights[b-1][a-1] = c;
    pair<int, VI> res = GetMinCut(weights);
    cout << "Case #" << i+1 << ": " << res.first << endl;
// END CUT
```

1.7 Graph cut inference

```
// Special-purpose \{0,1\} combinatorial optimization solver for
// problems of the following by a reduction to graph cuts:
                         sum_i psi_i(x[i]) + sum_{\{i < j\}} phi_{\{ij\}}(x[i], x[j])
         minimize
// x[1]...x[n] in {0,1}
       psi_i : {0, 1} --> R
   phi_{ij}: {0, 1} x {0, 1} --> R
// such that
   phi_{ij}(0,0) + phi_{ij}(1,1) \le phi_{ij}(0,1) + phi_{ij}(1,0)  (*)
// direction of the inequality in (\star) is reversed.
// INPUT: phi -- a matrix such that phi[i][j][u][v] = phi_{ij}(u, v)
         psi -- a matrix such that psi[i][u] = psi_i(u)
         x -- a vector where the optimal solution will be stored
// OUTPUT: value of the optimal solution
//\ {\it To\ use\ this\ code,\ create\ a\ GraphCutInference\ object,\ and\ call\ the}
// DoInference() method. To perform maximization instead of minimization,
// ensure that #define MAXIMIZATION is enabled.
#include <vector>
#include <iostream>
using namespace std:
typedef vector<int> VI;
typedef vector<VI> VVI;
typedef vector<VVI> VVVI;
```

```
typedef vector<VVVI> VVVVI;
const int INF = 1000000000;
// comment out following line for minimization
#define MAXIMIZATION
struct GraphCutInference {
  int N;
  VVI cap, flow;
  VI reached:
  int Augment(int s, int t, int a) {
    reached[s] = 1;
if (s == t) return a;
for (int k = 0; k < N; k++) {
   if (reached[k]) continue;</pre>
       if (int aa = min(a, cap[s][k] - flow[s][k])) {
          if (int b = Augment(k, t, aa)) {
            flow[s][k] += b;
             flow[k][s] -= b;
            return b;
     return 0:
  int GetMaxFlow(int s, int t) {
    N = cap.size();
     flow = VVI(N, VI(N));
     reached = VI(N);
     int totflow = 0;
     while (int amt = Augment(s, t, INF)) {
       totflow += amt;
       fill(reached.begin(), reached.end(), 0);
     return totflow;
  int DoInference(const VVVVI &phi, const VVI &psi, VI &x) {
    int M = phi.size():
     cap = VVI(M+2, VI(M+2));
     VI b(M);
     int c = 0;
     for (int i = 0; i < M; i++) {
  b[i] += psi[i][1] - psi[i][0];</pre>
       c += psi[i][0];
       C += psi[1][0];
for (int j = 0; j < i; j++)
  b[i] += phi[i][j][1][1] - phi[i][j][0][1];
for (int j = i+1; j < M; j++) {
  cap[i][j] = phi[i][j][0][1] + phi[i][j][1][0] - phi[i][j][0][0] - phi[i][j][1][1];
  b[i] += phi[i][j][1][0] - phi[i][j][0][0];
  c += phi[i][j][0][0];</pre>
#ifdef MAXIMIZATION
     for (int i = 0; i < M; i++) {
       for (int j = i+1; j < M; j++)
          cap[i][j] *= -1;
     c *= -1;
#endif
    for (int i = 0; i < M; i++) {
  if (b[i] >= 0) {
         cap[M][i] = b[i];
       } else {
         cap[i][M+1] = -b[i];
          c += b[i];
     int score = GetMaxFlow(M, M+1);
     fill(reached.begin(), reached.end(), 0);
     Augment (M, M+1, INF);
     x = VI(M);
     for (int i = 0; i < M; i++) x[i] = reached[i] ? 0 : 1;</pre>
      score += c:
#ifdef MAXIMIZATION
     score \star = -1:
#endif
     return score:
};
```

```
int main() {
  // solver for "Cat vs. Dog" from NWERC 2008
  cin >> numcases;
  for (int caseno = 0; caseno < numcases; caseno++) {</pre>
    int c, d, v;
    cin >> c >> d >> v;
    VVVVI phi(c+d, VVVI(c+d, VVI(2, VI(2))));
    VVI psi(c+d, VI(2));
for (int i = 0; i < v; i++) {</pre>
      char p, q;
      int u, v;
cin >> p >> u >> q >> v;
       u--; v--;
      if (p == 'C')
        phi[u][c+v][0][0]++;
         phi[c+v][u][0][0]++;
      } else {
        phi[v][c+u][1][1]++;
         phi[c+u][v][1][1]++;
    GraphCutInference graph;
    VI x;
    cout << graph.DoInference(phi, psi, x) << endl;</pre>
  return 0:
```

2 Geometry

2.1 Convex hull

```
// Compute the 2D convex hull of a set of points using the monotone chain
// algorithm. Eliminate redundant points from the hull if REMOVE_REDUNDANT is
// #defined.
// Running time: O(n log n)
    INPUT: a vector of input points, unordered.
    OUTPUT: a vector of points in the convex hull, counterclockwise, starting
              with bottommost/leftmost point
#include <cstdio>
#include <cassert>
#include <vector>
#include <algorithm>
#include <cmath>
 // BEGIN CUT
#include <map>
// END CUT
using namespace std;
#define REMOVE_REDUNDANT
typedef double T;
const T EPS = 1e-7;
struct PT (
  T x, y;
 PT() {}
 PT(T x, T y) : x(x), y(y) {}
bool operator<(const PT &rhs) const { return make_pair(y,x) < make_pair(rhs.y,rhs.x); }</pre>
 bool operator == (const PT &rhs) const { return make_pair(y,x) == make_pair(rhs.y,rhs.x); }
T cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
T area2(PT a, PT b, PT c) { return cross(a,b) + cross(b,c) + cross(c,a); }
#ifdef REMOVE_REDUNDANT
bool between (const PT &a, const PT &b, const PT &c) {
  return (fabs (area2 (a,b,c)) < EPS && (a.x-b.x)*(c.x-b.x) <= 0 && (a.y-b.y)*(c.y-b.y) <= 0);
#endif
void ConvexHull(vector<PT> &pts) {
 sort(pts.begin(), pts.end());
  pts.erase(unique(pts.begin(), pts.end()), pts.end());
  vector<PT> up, dn;
  for (int i = 0; i < pts.size(); i++) {</pre>
```

```
while (up.size() > 1 && area2(up[up.size()-2], up.back(), pts[i]) >= 0) up.pop_back();
    while (dn.size() > 1 && area2(dn[dn.size()-2], dn.back(), pts[i]) <= 0) dn.pop_back();
    up.push_back(pts[i]);
    dn.push_back(pts[i]);
  pts = dn;
  for (int i = (int) up.size() - 2; i >= 1; i--) pts.push_back(up[i]);
#ifdef REMOVE_REDUNDANT
  if (pts.size() <= 2) return;</pre>
  dn.clear();
  dn.push_back(pts[0]);
  dn.push back(pts[1]);
  for (int i = 2; i < pts.size(); i++) {
   if (between(dn[dn.size()-2], dn[dn.size()-1], pts[i])) dn.pop_back();</pre>
    dn.push_back(pts[i]);
  if (dn.size() >= 3 && between(dn.back(), dn[0], dn[1])) {
    dn[0] = dn.back();
    dn.pop_back();
#endif
// The following code solves SPOJ problem #26: Build the Fence (BSHEEP)
 int t:
  scanf("%d", &t);
  for (int caseno = 0; caseno < t; caseno++) {
    scanf("%d", &n);
    vector<PT> v(n);
    for (int i = 0; i < n; i++) scanf("%lf%lf", &v[i].x, &v[i].y);
    vector<PT> h(v);
    map<PT, int> index;
    for (int i = n-1; i >= 0; i--) index[v[i]] = i+1;
    ConvexHull(h):
    double len = 0;
    for (int i = 0; i < h.size(); i++) {
      double dx = h[i].x - h[(i+1)%h.size()].x;
      double dy = h[i].y - h[(i+1)%h.size()].y;
      len += sqrt (dx*dx+dy*dy);
    if (caseno > 0) printf("\n");
printf("%.2f\n", len);
    for (int i = 0; i < h.size(); i++) {
  if (i > 0) printf(" ");
  printf("%d", index[h[i]]);
    printf("\n");
// END CUT
```

2.2 Miscellaneous geometry

```
// C++ routines for computational geometry.
#include <iostream>
#include <vector>
#include <cmath>
#include <cassert>
using namespace std;
double INF = 1e100;
double EPS = 1e-12;
struct PT {
  double x, y;
  PT (double x, double y) : x(x), y(y) {}
  PT (const PT &p) : x(p.x), y(p.y)
  PT operator + (const PT &p) const { return PT(x+p.x, y+p.y); }
 PT operator - (const PT &p) const { return PT(x-p.x, y-p.y); }
  PT operator * (double c)
                              const { return PT(x*c, y*c ); ]
 PT operator / (double c)
                              const { return PT(x/c, y/c );
double dot(PT p, PT q)
                          { return p.x*q.x+p.y*q.y; }
double dist2(PT p, PT q) { return dot(p-q,p-q); }
```

```
double cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
ostream & operator << (ostream & os, const PT & p) {
    return os << "(" << p.x << "," << p.y << ")";
// rotate a point CCW or CW around the origin
PT RotateCCW90 (PT p) { return PT (-p.y,p.x);
                       { return PT(p.y,-p.x); }
PT RotateCW90(PT p)
PT RotateCCW(PT p, double t) {
  return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos(t));
// project point c onto line through a and b
// assuming a != b
PT ProjectPointLine(PT a, PT b, PT c) {
 return a + (b-a) *dot (c-a, b-a) /dot (b-a, b-a);
// project point c onto line segment through a and b
PT ProjectPointSegment(PT a, PT b, PT c) {
  double r = dot(b-a, b-a);
  if (fabs(r) < EPS) return a;</pre>
  r = dot(c-a, b-a)/r;
  if (r < 0) return a;</pre>
  if (r > 1) return b;
  return a + (b-a) *r;
// compute distance from c to segment between a and b
double DistancePointSegment (PT a, PT b, PT c) {
 return sqrt(dist2(c, ProjectPointSegment(a, b, c)));
// compute distance between point (x,y,z) and plane ax+by+cz=d
double DistancePointPlane(double x, double y, double z,
                          double a, double b, double c, double d)
  return fabs(a*x+b*y+c*z-d)/sqrt(a*a+b*b+c*c);
// determine if lines from a to b and c to d are parallel or collinear
bool LinesParallel(PT a, PT b, PT c, PT d) {
 return fabs(cross(b-a, c-d)) < EPS;
bool LinesCollinear(PT a, PT b, PT c, PT d) {
  return LinesParallel(a, b, c, d)
      && fabs(cross(a-b, a-c)) < EPS
      && fabs(cross(c-d, c-a)) < EPS;
// determine if line segment from a to b intersects with
// line seament from c to d
bool SegmentsIntersect(PT a, PT b, PT c, PT d) {
  dist2(b, c) < EPS || dist2(b, d) < EPS) return true;
    if (dot(c-a, c-b) > 0 && dot(d-a, d-b) > 0 && dot(c-b, d-b) > 0)
     return false;
  if (cross(d-a, b-a) * cross(c-a, b-a) > 0) return false;
  if (cross(a-c, d-c) * cross(b-c, d-c) > 0) return false;
// compute intersection of line passing through a and b
\ensuremath{//} with line passing through c and d, assuming that unique
// intersection exists; for segment intersection, check if
// segments intersect first
PT ComputeLineIntersection(PT a, PT b, PT c, PT d) {
  b=b-a; d=c-d; c=c-a;
  assert (dot (b, b) > EPS && dot (d, d) > EPS);
  return a + b*cross(c, d)/cross(b, d);
// compute center of circle given three points
PT ComputeCircleCenter(PT a, PT b, PT c) {
  c = (a+c)/2;
  return ComputeLineIntersection(b, b+RotateCW90(a-b), c, c+RotateCW90(a-c));
// determine if point is in a possibly non-convex polygon (by William
// Randolph Franklin); returns 1 for strictly interior points, 0 for
// strictly exterior points, and 0 or 1 for the remaining points.
// Note that it is possible to convert this into an *exact* test using
// integer arithmetic by taking care of the division appropriately
// (making sure to deal with signs properly) and then by writing exact
 // tests for checking point on polygon boundary
bool PointInPolygon(const vector<PT> &p, PT q) {
```

bool c = 0;

```
for (int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) %p.size();
    if ((p[i].y <= q.y && q.y < p[j].y ||
      p[j].y <= q.y && q.y < p[i].y) &&
       q.x < p[i].x + (p[j].x - p[i].x) * (q.y - p[i].y) / (p[j].y - p[i].y))
  return c;
// determine if point is on the boundary of a polygon
bool PointOnPolygon(const vector<PT> &p, PT q) {
  for (int i = 0; i < p.size(); i++)</pre>
     \textbf{if} \ (\texttt{dist2}(\texttt{ProjectPointSegment}(\texttt{p[i]},\ \texttt{p[(i+1)\$p.size()]},\ \texttt{q)},\ \texttt{q)} \ < \ \texttt{EPS)} 
      return true:
    return false:
// compute intersection of line through points a and b with
// circle centered at c with radius r > 0
vector<PT> CircleLineIntersection(PT a, PT b, PT c, double r) {
  vector<PT> ret:
  b = b-a:
  a = a-c:
  double A = dot(b, b);
  double B = dot(a, b);
  double C = dot(a, a) - r*r;
  double D = B*B - A*C:
  if (D < -EPS) return ret;</pre>
  ret.push back(c+a+b*(-B+sgrt(D+EPS))/A);
  if (D > EPS)
    ret.push_back(c+a+b*(-B-sqrt(D))/A);
  return ret;
// compute intersection of circle centered at a with radius r
// with circle centered at b with radius R
vector<PT> CircleCircleIntersection(PT a, PT b, double r, double R) {
  vector<PT> ret;
  double d = sqrt (dist2(a, b));
if (d > r+R || d+min(r, R) < max(r, R)) return ret;
double x = (d*d-R*R*r*r)/(2*d);</pre>
  double y = sqrt(r*r-x*x);
  PT v = (b-a)/d;
  ret.push back(a+v*x + RotateCCW90(v)*y);
  if (y > 0)
    ret.push_back(a+v*x - RotateCCW90(v)*y);
  return ret;
// This code computes the area or centroid of a (possibly nonconvex)
// polygon, assuming that the coordinates are listed in a clockwise or
// counterclockwise fashion. Note that the centroid is often known as // the "center of gravity" or "center of mass".
double ComputeSignedArea(const vector<PT> &p) {
  double area = 0;
  for(int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
    area += p[i].x*p[j].y - p[j].x*p[i].y;
  return area / 2.0;
double ComputeArea(const vector<PT> &p) {
  return fabs(ComputeSignedArea(p));
PT ComputeCentroid(const vector<PT> &p) {
  double scale = 6.0 * ComputeSignedArea(p);
  double scare = 0.0 * computerignedatea(p);
for (int i = 0; i < p.size();
  int j = (i+1) % p.size();
  c = c + (p[i]+p[j])*(p[i].x*p[j].y - p[j].x*p[i].y);</pre>
  return c / scale;
// tests whether or not a given polygon (in CW or CCW order) is simple
bool IsSimple(const vector<PT> &p) {
  for (int i = 0; i < p.size(); i++) {</pre>
    for (int k = i+1; k < p.size(); k++) {
      int j = (i+1) % p.size();
int l = (k+1) % p.size();
if (i == l || j == k) continue;
       if (SegmentsIntersect(p[i], p[j], p[k], p[l]))
         return false;
  return true:
int main() {
```

```
// expected: (-5,2)
cerr << RotateCCW90(PT(2,5)) << endl;
// expected: (5,-2)
cerr << RotateCW90(PT(2,5)) << endl;
// expected: (-5,2)
cerr << RotateCCW(PT(2,5),M_PI/2) << endl;</pre>
// expected: (5,2)
cerr << ProjectPointLine(PT(-5,-2), PT(10,4), PT(3,7)) << endl;</pre>
// expected: (5,2) (7.5,3) (2.5,1)
<< ProjectPointSegment(PT(-5,-2), PT(2.5,1), PT(3,7)) << endl;
// expected: 6.78903
cerr << DistancePointPlane(4,-4,3,2,-2,5,-8) << endl;</pre>
// expected: 1 0 1
cerr << LinesParallel(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "</pre>
      << LinesParallel(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << " "
      << LinesParallel(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;
// expected: 0 0 1
// expected: 1 1 1 0
cerr << SegmentsIntersect(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << " "
      << SegmentsIntersect(PT(0,0), PT(2,4), PT(4,3), PT(0,5)) << " "
      << SegmentsIntersect(PT(0,0), PT(2,4), PT(2,-1), PT(-2,1)) << " "
      << SegmentsIntersect(PT(0,0), PT(2,4), PT(5,5), PT(1,7)) << endl;
// expected: (1,2)
cerr << ComputeLineIntersection(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << endl;</pre>
// expected: (1,1)
cerr << ComputeCircleCenter(PT(-3,4), PT(6,1), PT(4,5)) << endl;</pre>
vector<PT> v;
v.push back(PT(0,0));
v.push_back(PT(5,0));
v.push_back(PT(5,5));
v.push_back(PT(0,5));
// expected: 1 1 1 0 0
cerr << PointInPolygon(v, PT(2,2)) << " "
      << PointInPolygon(v, PT(2,0)) << " "
      << PointInPolygon(v, PT(0,2)) << " "</pre>
      << PointInPolygon(v, PT(5,2)) << " "
      << PointInPolygon(v, PT(2,5)) << endl;
// expected: 0 1 1 1 1
cerr << PointOnPolygon(v, PT(2,2)) << " "</pre>
      << PointOnPolygon(v, PT(2,0)) << " "
      << PointOnPolygon(v, PT(0,2)) << " "
      << PointOnPolygon(v, PT(5,2)) << " "
      << PointOnPolygon(v, PT(2,5)) << endl;
               (5,4) (4,5)
               blank line
               (4,5) (5,4)
               blank line
               (4,5) (5,4)
vector<PT> u = CircleLineIntersection(PT(0,6), PT(2,6), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
u = CircleLineIntersection(PT(0,9), PT(9,0), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
u = CircleCircleIntersection(PT(1,1), PT(10,10), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
u = CircleCircleIntersection(PT(1,1), PT(8,8), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 10, sqrt(2.0)/2.0);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 5, sqrt(2.0)/2.0);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
// area should be 5.0
// centroid should be (1.1666666, 1.166666)
PT pa[] = \{ PT(0,0), PT(5,0), PT(1,1), PT(0,5) \};
vector<PT> p(pa, pa+4);
PT c = ComputeCentroid(p);
cerr << "Area: " << ComputeArea(p) << endl;
cerr << "Centroid: " << c << endl;
return 0:
```

3 Numerical algorithms

3.1 Number theory (modular, Chinese remainder, linear Diophantine)

```
// This is a collection of useful code for solving problems that
// involve modular linear equations. Note that all of the
// algorithms described here work on nonnegative integers.
#include <iostream>
#include <vector>
#include <algorithm>
using namespace std;
typedef vector<int> VI;
typedef pair<int, int> PII;
// return a % b (positive value)
int mod(int a, int b) {
        return ((a%b) + b) % b;
// computes gcd(a,b)
int gcd(int a, int b) {
    while (b) { int t = a%b; a = b; b = t; }
        return a;
// computes lcm(a,b)
int lcm(int a, int b)
        return a / gcd(a, b) *b;
// (a^b) mod m via successive squaring
int powermod(int a, int b, int m)
        int ret = 1;
        while (b)
                if (b & 1) ret = mod(ret*a, m);
                a = mod(a*a, m);
// returns g = gcd(a, b); finds x, y such that d = ax + by
int extended_euclid(int a, int b, int &x, int &y) {
        int xx = y = 0;
        int yy = x = 1;
        while (b) {
                int q = a / b;
                int t = b; b = a%b; a = t;
                t = xx; xx = x - q*xx; x = t;
                t = yy; yy = y - q*yy; y = t;
// finds all solutions to ax = b \pmod{n}
VI modular_linear_equation_solver(int a, int b, int n) {
        int x, y;
        VI ret:
        int g = extended_euclid(a, n, x, y);
        if (!(b%g)) {
                x = mod(x*(b / g), n);
                for (int i = 0; i < q; i++)
                        ret.push_back(mod(x + i*(n / g), n));
// computes b such that ab = 1 \pmod{n}, returns -1 on failure
int mod_inverse(int a, int n) {
        int g = extended_euclid(a, n, x, y);
        if (g > 1) return -1;
        return mod(x, n);
// Chinese remainder theorem (special case): find z such that
// z % m1 = r1, z % m2 = r2. Here, z is unique modulo M = lcm(m1, m2).
// Return (z, M). On failure, M = -1.
```

```
PII chinese_remainder_theorem(int m1, int r1, int m2, int r2) {
         int s, t;
         int g = extended_euclid(m1, m2, s, t);
         if (r1%g != r2%g) return make_pair(0, -1);
         return make_pair(mod(s*r2*m1 + t*r1*m2, m1*m2) / g, m1*m2 / g);
// Chinese remainder theorem: find z such that
// z % m[i] = r[i] for all i. Note that the solution is
// unique modulo M = lcm_i (m[i]). Return (z, M). On
// failure, M = -1. Note that we do not require the a[i]'s
// to be relatively prime.
PII chinese_remainder_theorem(const VI &m, const VI &r) {
         PII ret = make_pair(r[0], m[0]);
for (int i = 1; i < m.size(); i++) {
    ret = chinese_remainder_theorem(ret.second, ret.first, m[i], r[i]);
}</pre>
                   if (ret.second == -1) break;
         return ret;
// computes x and y such that ax + by = c
// returns whether the solution exists
bool linear_diophantine(int a, int b, int c, int &x, int &y) {
         if (!a && !b)
                   if (c) return false:
                   x = 0; v = 0;
                  return true;
         if (!a)
                   if (c % b) return false;
                   return true;
         if (!b)
                  if (c % a) return false;
x = c / a; y = 0;
                  return true;
         int q = \gcd(a, b);
         if (c % g) return false;
         x = c / g * mod_inverse(a / g, b / g);
         y = (c - a*x) / b;
         return true;
int main() {
         // expected: 2
         cout << gcd(14, 30) << endl;
         // expected: 2 -2 1
         int x, y;
int g = extended_euclid(14, 30, x, y);
cout << g << " " << x << " " << y << endl;</pre>
         // expected: 95 451
         VI sols = modular_linear_equation_solver(14, 30, 100);
         for (int i = 0; i < sols.size(); i++) cout << sols[i] << " ";</pre>
         cout << mod_inverse(8, 9) << endl;</pre>
         // expected: 23 105
                      11 12
         PII ret = chinese_remainder_theorem(VI({ 3, 5, 7 }), VI({ 2, 3, 2 })); cout << ret.first << " " << ret.second << endl;
         ret = chinese_remainder_theorem(VI({ 4, 6 }), VI({ 3, 5 }));
cout << ret.first << " " << ret.second << end;
            expected: 5 -15
         if (!linear_diophantine(7, 2, 5, x, y)) cout << "ERROR" << endl;
cout << x << " " << y << endl;</pre>
         return 0;
```

3.2 Systems of linear equations, matrix inverse, determinant

```
// Gauss-Jordan elimination with full pivoting.
//
// Uses:
// (1) solving systems of linear equations (AX=B)
```

```
(2) inverting matrices (AX=I)
      (3) computing determinants of square matrices
// Running time: O(n^3)
                  a[][] = an nxn matrix
                  b[][] = an nxm matrix
// OUTPUT: X
                         = an nxm matrix (stored in b[][])
                  A^{-1} = an nxn matrix (stored in a[][])
                  returns determinant of a[][]
#include <iostream>
#include <vector>
#include <cmath>
using namespace std;
const double EPS = 1e-10;
typedef vector<int> VI;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
T GaussJordan (VVT &a, VVT &b) {
  const int n = a.size();
   const int m = b[0].size();
   VI irow(n), icol(n), ipiv(n);
  T det = 1:
   for (int i = 0; i < n; i++) {</pre>
     for (int j = 0, j < n, j++) {
  int pj = -1, pk = -1;
  for (int j = 0; j < n; j++) if (!ipiv[j])
    for (int k = 0; k < n; k++) if (!ipiv[k])</pre>
     if (pj == -1 || fabs(a[j][k]) > fabs(a[pj][pk])) { pj = j; pk = k; }
if (fabs(a[pj][pk]) < EPS) { cerr << "Matrix is singular." << endl; exit(0); }</pre>
     ipiv[pk]++;
     swap(a[pj], a[pk]);
     swap(b[pj], b[pk]);
     if (pj != pk) det *= -1;
irow[i] = pj;
icol[i] = pk;
     T c = 1.0 / a[pk][pk];
     det *= a[pk][pk];
     a[pk][pk] = 1.0;
     for (int p = 0; p < n; p++) a[pk][p] *= c;
for (int p = 0; p < m; p++) b[pk][p] *= c;
for (int p = 0; p < n; p++) if (p != pk) {
        c = a[p][pk];
        for (int q = 0; q < n; q++) a[p][q] -= a[pk][q] * c;
for (int q = 0; q < m; q++) b[p][q] -= b[pk][q] * c;
  for (int p = n-1; p >= 0; p--) if (irow[p] != icol[p]) {
   for (int k = 0; k < n; k++) swap(a[k][irow[p]], a[k][icol[p]]);</pre>
   return det;
int main() {
   const int n = 4;
   const int m = 2;
  double A[n][n] = \{ \{1,2,3,4\}, \{1,0,1,0\}, \{5,3,2,4\}, \{6,1,4,6\} \}; double B[n][m] = \{ \{1,2\}, \{4,3\}, \{5,6\}, \{8,7\} \};
   VVT a(n), b(n);
  for (int i = 0; i < n; i++) {
   a[i] = VT(A[i], A[i] + n);
   b[i] = VT(B[i], B[i] + m);</pre>
   double det = GaussJordan(a, b);
   // expected: 60
   cout << "Determinant: " << det << endl;</pre>
   // expected: -0.233333 0.166667 0.133333 0.0666667
                     0.166667 0.166667 0.333333 -0.333333 
0.233333 0.833333 -0.133333 -0.0666667
                     0.05 -0.75 -0.1 0.2
   cout << "Inverse: " << endl;</pre>
   for (int i = 0; i < n; i++) {
  for (int j = 0; j < n; j++)
    cout << a[i][j] << ' ';</pre>
     cout << endl;
```

// expected: 1.63333 1.3

```
// -0.166667 0.5

// 2.36667 1.7

// 2.36667 1.7

// 3.3667 1.7

cout << "Solution: " << endl;

for (int i = 0; i < n; i++) {

   for (int j = 0; j < m; j++) {

      cout << b[i][j] << ' ';

   cout << endl;

}
```

3.3 Reduced row echelon form, matrix rank

```
// Reduced row echelon form via Gauss-Jordan elimination
// with partial pivoting. This can be used for computing
// the rank of a matrix.
// Running time: O(n^3)
// INPUT: a[][] = an nxm matrix
// OUTPUT: rref[][] = an nxm matrix (stored in a[][])
              returns rank of a[][]
#include <iostream>
#include <vector>
#include <cmath>
using namespace std;
const double EPSILON = 1e-10;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
int rref(VVT &a) {
 int n = a.size();
int m = a[0].size();
  int r = 0;
  for (int c = 0; c < m && r < n; c++) {
    int i = r:
    for (int i = r + 1; i < n; i++)
      if (fabs(a[i][c]) > fabs(a[j][c])) j = i;
    if (fabs(a[j][c]) < EPSILON) continue;</pre>
    swap(a[j], a[r]);
    T s = 1.0 / a[r][c];
    for (int j = 0; j < m; j++) a[r][j] *= s;
for (int i = 0; i < n; i++) if (i != r) {</pre>
      T t = a[i][c];
      for (int j = 0; j < m; j++) a[i][j] -= t * a[r][j];</pre>
    r++;
  return r;
int main() {
  const int n = 5, m = 4;
  double A[n][m] = {
    {16, 2, 3, 13},
    { 5, 11, 10, 8},
    { 9, 7, 6, 12}, { 4, 14, 15, 1},
    {13, 21, 21, 13}};
  VVT a(n);
  for (int i = 0; i < n; i++)
    a[i] = VT(A[i], A[i] + m);
  int rank = rref(a);
  // expected: 3
  cout << "Rank: " << rank << endl;
  // expected: 1 0 0 1
              0 1 0 3
                0 0 1 -3
                0 0 0 3.10862e-15
                0 0 0 2.22045e-15
  cout << "rref: " << endl;
  for (int i = 0; i < 5; i++) {
  for (int j = 0; j < 4; j++)
    cout << a[i][j] << ' ';
    cout << endl;
```

3.4 Fast Fourier transform

```
#include <cassert>
#include <cstdio>
#include <cmath>
struct cpx
  cpx(){}
  cpx (double aa):a(aa),b(0){}
  cpx(double aa, double bb):a(aa),b(bb){}
  double a;
  double b:
  double modsq(void) const
    return a * a + b * b;
  cpx bar(void) const
    return cpx(a, -b);
cpx operator + (cpx a, cpx b)
  return cpx(a.a + b.a, a.b + b.b);
cpx operator * (cpx a, cpx b)
  return cpx(a.a * b.a - a.b * b.b, a.a * b.b + a.b * b.a);
cpx operator / (cpx a, cpx b)
  cpx r = a * b.bar();
  return cpx(r.a / b.modsq(), r.b / b.modsq());
cpx EXP (double theta)
  return cpx(cos(theta),sin(theta));
const double two_pi = 4 * acos(0);
// in:
           input array
// out:
          output array
// step: {SET TO 1} (used internally)
// size: length of the input/output {MUST BE A POWER OF 2}
          either plus or minus one (direction of the FFT)
// RESULT: out[k] = \sum_{j=0}^{size} - 1} in[j] * exp(dir * 2pi * i * j * k / size)
void FFT(cpx *in, cpx *out, int step, int size, int dir)
  if(size < 1) return;</pre>
  if(size == 1)
    out[0] = in[0];
    return;
  FFT(in, out, step * 2, size / 2, dir);
  FFT(in + step, out + size / 2, step * 2, size / 2, dir);
  for (int i = 0; i < size / 2; i++)
    cpx even = out[i];
    cpx odd = out[i + size / 2];
    out[i] = even + EXP(dir * two_pi * i / size) * odd;
out[i + size / 2] = even + EXP(dir * two_pi * (i + size / 2) / size) * odd;
// Usage:
// f[0...N-1] and q[0..N-1] are numbers
// Want to compute the convolution h, defined by
// h[n] = sum of f[k]g[n-k] (k = 0, ..., N-1).

// Here, the index is cyclic; f[-1] = f[N-1], f[-2] = f[N-2], etc.
// Let F[0...N-1] be FFT(f), and similarly, define G and H.
// The convolution theorem says H[n] = F[n]G[n] (element-wise product).
// To compute h[] in O(N log N) time, do the following:
    1. Compute F and G (pass dir = 1 as the argument).
// 2. Get H by element-wise multiplying F and G.
    3. Get h by taking the inverse FFT (use dir = -1 as the argument)
        and *dividing by N*. DO NOT FORGET THIS SCALING FACTOR.
int main (void)
  printf("If rows come in identical pairs, then everything works.\n");
  cpx \ a[8] = \{0, 1, cpx(1,3), cpx(0,5), 1, 0, 2, 0\};
```

```
cpx b[8] = \{1, cpx(0,-2), cpx(0,1), 3, -1, -3, 1, -2\};
cpx A[8];
cpx B[8];
FFT(a, A, 1, 8, 1);
FFT(b, B, 1, 8, 1);
for (int i = 0; i < 8; i++)
  printf("%7.21f%7.21f", A[i].a, A[i].b);
printf("\n");
for(int i = 0; i < 8; i++)
  cpx Ai(0,0);
  for (int j = 0; j < 8; j++)
    Ai = Ai + a[j] * EXP(j * i * two_pi / 8);
  printf("%7.21f%7.21f", Ai.a, Ai.b);
printf("\n");
cpx AB[8];
for(int i = 0 ; i < 8 ; i++)</pre>
 AB[i] = A[i] * B[i];
cpx aconvb[8];
FFT (AB, aconvb, 1, 8, -1);
for(int i = 0 ; i < 8 ; i++)
aconvb[i] = aconvb[i] / 8;
for(int i = 0; i < 8; i++)
  printf("%7.21f%7.21f", aconvb[i].a, aconvb[i].b);
printf("\n");
for(int i = 0 ; i < 8 ; i++)
  cpx aconvbi(0,0);
  for (int j = 0; j < 8; j++)
    aconvbi = aconvbi + a[j] * b[(8 + i - j) % 8];
  printf("%7.21f%7.21f", aconvbi.a, aconvbi.b);
printf("\n");
return 0;
```

3.5 Simplex algorithm

```
// Two-phase simplex algorithm for solving linear programs of the form
       maximize
       subject to Ax <= b
                   x >= 0
// INPUT: A -- an m x n matrix
         b -- an m-dimensional vector
         c -- an n-dimensional vector
          x -- a vector where the optimal solution will be stored
// OUTPUT: value of the optimal solution (infinity if unbounded
          above, nan if infeasible)
// To use this code, create an LPSolver object with A, b, and c as
// arguments. Then, call Solve(x).
#include <iostream>
#include <iomanip>
#include <vector>
#include <cmath>
#include <limits>
using namespace std;
typedef long double DOUBLE;
typedef vector<DOUBLE> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
const DOUBLE EPS = 1e-9;
struct LPSolver {
 int m, n;
  VI B. N.
  VVD D;
```

```
LPSolver(const VVD &A, const VD &b, const VD &c) :
      m(b.size()), n(c.size()), N(n + 1), B(m), D(m + 2, VD(n + 2)) {
      for (int i = 0; i < m; i++) for (int j = 0; j < n; j++) D[i][j] = A[i][j];
for (int i = 0; i < m; i++) { B[i] = n + i; D[i][n] = -1; D[i][n + 1] = b[i]; }</pre>
      for (int j = 0; j < n; j++) { N[j] = j; D[m][j] = -c[j]; }</pre>
      N[n] = -1; D[m + 1][n] = 1;
   void Pivot(int r, int s)
     double inv = 1.0 / D[r][s];
for (int i = 0; i < m + 2; i++) if (i != r)
    for (int j = 0; j < n + 2; j++) if (j != s)
        D[i][j] -= D[r][j] * D[i][s] * inv;
    for (int j = 0; j < n + 2; j++) if (j != s) D[r][j] *= inv;
    for (int i = 0; i < m + 2; i++) if (i != r) D[i][s] *= -inv;
    for (int i = 0; i < m + 2; i++) if (i != r) D[i][s] *= -inv;</pre>
      D[r][s] = inv;
      swap(B[r], N[s]);
   bool Simplex(int phase) {
      int x = phase == 1 ? m + 1 : m;
      while (true) {
        int s = -1;
        for (int j = 0; j <= n; j++) {
   if (phase == 2 && N[j] == -1) continue;
   if (s == -1 || D[x][j] < D[x][s] || D[x][j] == D[x][s] && N[j] < N[s]) s = j;</pre>
         if (D[x][s] > -EPS) return true;
         int r = -1;
         for (int i = 0; i < m; i++) {
           if (D[i][s] < EPS) continue;
if (r == -1 || D[i][n + 1] / D[i][s] < D[r][n + 1] / D[r][s] ||
(D[i][n + 1] / D[i][s]) == (D[r][n + 1] / D[r][s]) && B[i] < B[r]) r = i;</pre>
         if (r == -1) return false;
        Pivot(r, s);
   DOUBLE Solve(VD &x) {
     int r = 0;
     for (int i = 1; i < m; i++) if (D[i][n + 1] < D[r][n + 1]) r = i;
if (D[r][n + 1] < -EPS) {</pre>
         Pivot(r, n);
         if (!Simplex(1) || D[m + 1][n + 1] < -EPS) return -numeric limits<DOUBLE>::infinity();
         for (int i = 0; i < m; i++) if (B[i] == -1) {
           int s = -1;
           for (int j = 0; j <= n; j++)

if (s == -1 \mid \mid D[i][j] < D[i][s] \mid \mid D[i][j] == D[i][s] && N[j] < N[s]) s = j;
      if (!Simplex(2)) return numeric_limits<DOUBLE>::infinity();
      x = VD(n);
      for (int i = 0; i < m; i++) if (B[i] < n) x[B[i]] = D[i][n + 1];
      return D[m][n + 1];
};
int main() {
   const int n = 3;
   DOUBLE _A[m][n] =
      { 6, -1, 0 },
      \{-1, -5, 0\},
      { 1, 5, 1 },
      { -1, -5, -1 }
   DOUBLE _b[m] = { 10, -4, 5, -5 };

DOUBLE _c[n] = { 1, -1, 0 };
   VD b(_b, _b + m);
VD c(_c, _c + n);
   for (int i = 0; i < m; i++) A[i] = VD(_A[i], _A[i] + n);</pre>
   LPSolver solver (A, b, c);
   DOUBLE value = solver.Solve(x);
   cerr << "VALUE: " << value << endl; // VALUE: 1.29032
cerr << "SOLUTION:"; // SOLUTION: 1.74194 0.451613 1
for (size_t i = 0; i < x.size(); i++) cerr << " " << x[i];</pre>
   cerr << endl:
   return 0:
```

4 Graph algorithms

4.1 Floyd's algorithm (C++)

```
#include <bits/stdc++.h>
using namespace std;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
typedef vector<int> VI;
typedef vector<VI> VVI;
// This function runs the Floyd-Warshall algorithm for all-pairs
// shortest paths. Also handles negative edge weights. Returns true
// if a negative weight cycle is found.
// Running time: O(|V|^3)
     INPUT: w[i][j] = weight of edge from i to j
OUTPUT: w[i][j] = shortest path from i to j
              prev[i][j] = node before j on the best path starting at i
bool FloydWarshall (VVT &w, VVI &prev) {
    int n = w.size();
    prev = VVI (n, VI(n, -1));
    for (int k = 0; k < n; k++) {
         for (int i = 0; i < n; i++) {</pre>
             for (int j = 0; j < n; j++) {
   if (w[i][j] > w[i][k] + w[k][j]) {
                      w[i][j] = w[i][k] + w[k][j];
                      prev[i][j] = k;
     // check for negative weight cycles
    for (int i=0;i<n;i++)</pre>
        if (w[i][i] < 0) return false;</pre>
    return true;
```

4.2 Topological sort (C++)

```
// This function uses performs a non-recursive topological sort.
// Running time: O(|V|^2). If you use adjacency lists (vector<map<int> >),
                 the running time is reduced to O(|E|).
    INPUT: w[i][j] = 1 if i should come before j, 0 otherwise
    OUTPUT: a permutation of 0,...,n-1 (stored in a vector)
              which represents an ordering of the nodes which
              is consistent with w
// If no ordering is possible, false is returned.
#include <iostream>
#include <queue>
#include <cmath>
#include <vector>
using namespace std;
typedef vector<T> VT;
typedef vector<VT> VVT;
typedef vector<int> VI;
typedef vector<VI> VVI;
bool TopologicalSort (const VVI &w, VI &order) {
 int n = w.size();
VI parents (n);
  queue<int> q;
 order.clear();
  for (int i = 0; i < n; i++) {
```

```
for (int j = 0; j < n; j++)
    if (w[j][i]) parents[i]++;
    if (parents[i] == 0) q.push (i);
}
while (q.size() > 0){
    int i = q.front();
    q.pop();
    order.push_back (i);
    for (int j = 0; j < n; j++) if (w[i][j]){
        parents[j]--;
        if (parents[j] == 0) q.push (j);
    }
}
return (order.size() == n);</pre>
```

4.3 Dijkstra's algorithm

```
#include <bits/stdc++.h>
using namespace std;
// (u)int64, float
template<typename D>
struct Dijkstra {
    static constexpr D Inf = 11 << 60;</pre>
    struct Edge { size_t to; D len; };
    struct Vertex {
        vector<Edge> outs;
#ifdef DIJKSTRA_RECORD_ROUTE
        size_t prev = -1;
#endif
    };
    size_t N;
    vector<Vertex> vs:
    // n nodes
    Dijkstra(size_t n) : N(n) {
        vs.resize(n):
    void add(size_t from, size_t to, D len) {
        assert (from < N);
        assert (to < N);
        assert (len >= 0);
        vs[from].outs.push_back({ to, len });
    void add_u(size_t a, size_t b, D len) {
        add(a, b, len);
        add(b, a, len);
    D solve(size_t from, size_t to) {
         vs[from].dist = 0;
#ifdef DIJKSTRA_RECORD_ROUTE
        vs[from].prev = from;
#endif
        auto comp = [&](size_t x, size_t y) {
    return vs[x].dist < vs[y].dist || (vs[x].dist == vs[y].dist && x < y);</pre>
        set<size_t, decltype(comp) > q { comp };
for (size_t i = 0; i < N; ++i) {</pre>
            q.insert(i);
        while (!q.empty()) {
            size_t i;
                 auto it = q.begin();
                 q.erase(it);
             if (i == to) {
                 goto RETURN;
             for (Edge const& e : vs[i].outs) {
                 if (vs[e.to].dist > vs[i].dist + e.len) {
                     if (q.find(e.to) != q.end()) {
```

4.4 SPFA shortest paths

```
// Shortest Path
template <long N>
struct SPFA {
               struct Edge {
                            long to, len;
               vector<Edge> edges;
               vector<long> outs[N];
               \begin{tabular}{ll} \beg
                            edges.push_back({to, len});
outs[from].push_back(edges.size() - 1);
               long dist[N];
               long route[N];
               long visiting[N];
               bool active[N];
               long solve(long from, long to) {
                             memset(active, 0, sizeof(active));
                             for (int i = 0; i < N; ++i) {
                                            dist[i] = 11 << 60;
                             long head = 0;
                            long tail = 0;
                             dist[from] = 0;
                             visiting[(tail++) % N] = from;
                             active[from] = true;
                             while (head < tail) {</pre>
                                           long i = visiting[(head++) % N];
active[i] = false;
                                            for (long j = 0; j < outs[i].size(); ++j) {</pre>
                                                          Edge &e = edges[outs[i][j]];
                                                           if (dist[e.to] > dist[i] + e.len) {
                                                                         dist[e.to] = dist[i] + e.len;
                                                                         route[e.to] = i;
                                                                         if (!active[e.to]) {
                                                                                        visiting[(tail++) % N] = e.to;
                                                                                        active[e.to] = true;
                             return dist[to];
};
```

4.5 Minimum spanning trees using Prim

```
#include <bits/stdc++.h>
using namespace std;

// (u)int64, float
template<typename D>
```

```
struct Prim {
    static constexpr D Inf = 11 << 60;</pre>
    struct Edge { size_t to; D len; };
        vector<Edge> outs;
        D dist = Inf;
        size_t prev = -1;
    size_t N;
    vector<Vertex> vs;
     // n nodes
    Prim(size_t n) : N(n) {
        vs.resize(n);
    void add(size_t from, size_t to, D len) {
        assert (from < N);
        assert (to < N);
        assert(len >= 0);
        vs[from].outs.push_back({ to, len });
    void add_u(size_t a, size_t b, D len) {
        add(a, b, len);
        add(b, a, len);
    D solve(size_t from) {
        vs[from].dist = 0;
        vs[from] prev = from;
        auto comp = [&](size_t x, size_t y) {
            return vs[x].dist < vs[y].dist || (vs[x].dist == vs[y].dist && x < y);
        set<size_t, decltype(comp)> q { comp };
        for (size_t i = 0; i < N; ++i) {</pre>
            q.insert(i);
        while (!q.empty()) {
            size_t i;
                auto it = q.begin();
                         = *it;
                q.erase(it);
            if (vs[i].dist == Inf) {
                vs[i].dist = 0;
            for (Edge const& e : vs[i].outs) {
                if (vs[e.to].dist > e.len) {
                    if (q.find(e.to) != q.end()) {
                         q.erase(e.to);
                         vs[e.to].dist = e.len;
                         vs[e.to].prev = i;
                         q.insert(e.to);
        D result = 0;
        for (size_t i = 0; i < N; ++i)
    result += vs[i].dist;</pre>
        return result;
};
```

4.6 Minimum spanning trees using Kruskal

```
#include <algorithm>
#include <cassert>
#include <cstdint>
#include <set>
#include <set>
#include <unordered_map>
#include <vector>

using namespace std;

// (u) int64, float
template <typename D>
```

```
struct Kruskal {
    static constexpr D Inf = 11 << 60;
    struct Edge {
        size_t n1, n2;
        bool operator<(Edge const& o) const noexcept {</pre>
            return len < o.len || (len == o.len && (n1 < o.n1 || (n1 == o.n1 && (n2 < o.n2))));
    typedef unordered_map<size_t, vector<Edge>> Forest;
    size_t N;
    vector<Edge> edges;
    Kruskal(size_t n) : N(n) {}
    void add(size_t n1, size_t n2, D len) {
        assert (n1 < N);
        assert (n2 < N);
        assert(len >= 0);
        if (n1 > n2) swap(n1, n2);
        edges.push_back({ n1, n2, len });
    void solve(Forest& forest) {
        sort(edges.begin(), edges.end());
        vector<size_t> _ufs(N);
for (size_t i = 0; i < N; i++)
   _ufs[i] = i;</pre>
        auto ufs_p = [&](size_t i) -> size_t* {
            auto p = i;
            while (_ufs[p] != p)
             p = _ufs[p];
_ufs[i] = p;
            return &_ufs[i];
        for (Edge const& e : edges) {
            auto* ufs_n1p = ufs_p(e.n1);
            auto* ufs_n2p = ufs_p(e.n2);
            if (*ufs_n1p != *ufs_n2p) {
                if (*ufs_nlp > *ufs_n2p)
                     swap(*ufs_n1p, *ufs_n2p);
                 auto it = forest.find(*ufs_n2p);
                 if (it != forest.end()) {
                     forest[*ufs_nlp].insert(forest[*ufs_nlp].end(), it->second.begin(), it->second.end
                           ());
                     forest.erase(it);
                 forest[*ufs_nlp].push_back(e);
                 *ufs_n2p = *ufs_n1p;
        return;
};
```

4.7 Strongly connected components

```
// Strongly Connected Components

template <long N>
struct Tarjan {
    vector<long> outs[N];

    void add(long from, long to) {
        outs[from].push_back(to);
    }

    long id_self[N];
    long id_low[N];
    long route[N];

    void dfs(long (&sec)[N], long from, long &last, long &now, long &now_sec) {
        id_self[from] = now;
        id_low[from] = now;
        now += 1;
```

```
route[from] = last;
    for (long j = 0; j < outs[from].size(); ++j) {
   long to = outs[from][j];</pre>
         if (!id_self[to]) {
             dfs(scc, to, last, now, now_scc);
             id_low[from] = min(id_low[from], id_low[to]);
         } else if (!scc[to]) {
             id_low[from] = min(id_low[from], id_self[to]);
    if (id_low[from] == id_self[from]) {
         while (last != from) {
             scc[last] = now_scc;
             last = route[last];
         scc[last] = now_scc;
         last = route[last];
         now_scc += 1;
void solve(long (&scc)[N]) {
    memset(id_self, 0, sizeof(id_self));
    long last = 0;
    long now_scc = 1;
    for (long i = 0; i < N; ++i) {
        if (!id_self[i]) {
             dfs(scc, i, last, now, now_scc);
    for (long i = 0; i < N; ++i) {
    scc[i] -= 1;</pre>
```

4.8 Eulerian path

};

```
// Eulerian Circuit
template <long N>
struct HierholzerUndirected {
   struct Edge {
       long to;
        bool chosen;
    vector<Edge> edges;
    vector<long> outs[N];
    void add(long from, long to) {
        edges.push_back(Edge {to, false});
        outs[from].push_back(edges.size() - 1);
        edges.push_back(Edge {from, false});
        outs[to].push_back(edges.size() - 1);
    long step[N];
    void dfs(vector<long> &path, long from) {
        for (; step[from] < outs[from].size(); ++step[from]) {</pre>
            Edge &e = edges[outs[from][step[from]]];
            if (!e.chosen) {
                edges[outs[from][step[from]] ^ 1].chosen = true;
                dfs(path, e.to);
        path.push_back(from);
    void solve(vector<long> &path, long from) {
        memset(step, 0, sizeof(step));
```

```
dfs(path, from);
};
template <long N>
struct HierholzerDirected {
    struct Edge {
        long to;
        bool chosen;
    };
    vector<Edge> edges;
    vector<long> outs[N];
    void add(long from, long to) {
   edges.push_back(Edge {to, false});
        outs[from].push_back(edges.size() - 1);
    long step[N];
    void dfs(vector<long> &path, long from) {
        for (; step[from] < outs[from].size(); ++step[from]) {</pre>
            Edge &e = edges[outs[from][step[from]]];
            if (!e.chosen) {
                 e.chosen = true;
                 dfs(path, e.to);
        path.push_back(from);
    void solve(vector<long> &path, long from) {
        memset(step, 0, sizeof(step));
        dfs(path, from);
};
```

4.9 Travelling Salesman Problem

```
// Travelling Salesman Problem
template <long N>
struct HeldKarp {
                 struct Edge
                                 long to, len;
                 vector<Edge> edges;
                 vector<long> outs[N];
                 \begin{tabular}{ll} \beg
                                  edges.push_back({to, len});
outs[from].push_back(edges.size() - 1);
                 long dist[N][N];
                 long best [11 << (N - 1)][N];
                 long solve(long n) {
                                  for (long i = 0; i < n; ++i) {
                                                   for (long j = 0; j < n; ++j) {
    dist[i][j] = 11 << 60;</pre>
                                                   dist[i][i] = 0;
                                                   for (long j = 0; j < outs[i].size(); ++j) {
    Edge &e = edges[outs[i][j]];</pre>
                                                                      dist[i][e.to] = min(dist[i][e.to], e.len);
                                  for (long i = 0; i < n; ++i) {
                                                   for (long j = 0; j < n; ++j) {
   for (long k = 0; k < n; ++k) {</pre>
                                                                                     dist[j][k] = min(dist[j][k], dist[j][i] + dist[i][k]);
                                  for (long i = 0; i < (11 << (n - 1)); ++i) {
                                                    for (long j = 0; j < n; ++j) {
```

```
best[i][j] = 11 << 60;
}

for (long j = 0; j < n - 1; ++j) {
    if (i == 11 << j) {
        best[i][j] = dist[n - 1][j];
    } else if (i & (11 << j)) {
        for (long k = 0; k < n - 1; ++k) {
            if ((i ^ (11 << j)) & (11 << k)) {
                best[i][j] = min(best[i][j], dist[j][k] + best[i ^ (11 << j))][k]);
            }
        }
    }
}

long result = 11 << 60;
for (long i = 0; i < n - 1; ++i) {
    result = min(result, dist[i][n - 1] + best[(11 << (n - 1)) - 1][i]);
}

return result;
}</pre>
```

5 Data structures

5.1 Suffix array

};

```
// Suffix array construction in O(L log^2 L) time. Routine for
// computing the length of the longest common prefix of any two
// suffixes in O(log L) time.
// INPUT: string s
// OUTPUT: array suffix[] such that suffix[i] = index (from 0 to L-1)
            of substring s[i...L-1] in the list of sorted suffixes.
            That is, if we take the inverse of the permutation suffix[],
            we get the actual suffix array.
#include <vector>
#include <iostream>
#include <string>
using namespace std;
struct SuffixArray {
  const int L;
  string s;
  vector<vector<int> > P:
  vector<pair<pair<int,int>,int> > M;
  SuffixArray(const string &s) : L(s.length()), s(s), P(1, vector<int>(L, 0)), M(L) {
   for (int i = 0; i < L; i++) P[0][i] = int(s[i]);
for (int skip = 1, level = 1; skip < L; skip *= 2, level++) {</pre>
      P.push_back(vector<int>(L, 0));
      for (int i = 0; i < L; i++)
       M[i] = make\_pair(make\_pair(P[level-1][i], i + skip < L ? P[level-1][i + skip] : -1000), i);
      sort (M.begin(), M.end());
         P[level][M[i].second] = (i > 0 \&\& M[i].first == M[i-1].first) ? P[level][M[i-1].second] : i; 
  vector<int> GetSuffixArray() { return P.back(); }
  // returns the length of the longest common prefix of s[i...L-1] and s[j...L-1]
  int LongestCommonPrefix(int i, int j) {
    int len = 0;
    if (i == j) return L - i;
    for (int k = P.size() - 1; k >= 0 && i < L && j < L; k--) {
     if (P[k][i] == P[k][j]) {
         j += 1 << k;
         len += 1 << k;
    return len:
};
// BEGIN CUT
   The following code solves UVA problem 11512: GATTACA.
#define TESTING
#ifdef TESTING
```

```
int main() {
 int T;
  for (int caseno = 0; caseno < T; caseno++) {</pre>
    cin >> s;
    SuffixArray array(s);
    vector<int> v = array.GetSuffixArray();
    int bestlen = -1, bestpos = -1, bestcount = 0;
    for (int i = 0; i < s.length(); i++) {</pre>
      int len = 0, count = 0;
for (int j = i+1; j < s.length(); j++) {</pre>
        int 1 = array.LongestCommonPrefix(i, j);
        if (1 >= len) {
         if (1 > len) count = 2; else count++;
          len = 1;
      if (len > bestlen || len == bestlen && s.substr(bestpos, bestlen) > s.substr(i, len)) {
        bestlen = len;
        bestcount = count;
        bestpos = i;
    if (bestlen == 0) {
      cout << "No repetitions found!" << endl;
    } else {
      cout << s.substr(bestpos, bestlen) << " " << bestcount << endl;</pre>
#else
// END CUT
int main() {
  // bobocel is the O'th suffix
  // obocel is the 5'th suffix
      bocel is the 1'st suffix
       ocel is the 6'th suffix
        cel is the 2'nd suffix
         el is the 3'rd suffix
           l is the 4'th suffix
  SuffixArray suffix("bobocel");
  vector<int> v = suffix.GetSuffixArray();
  // Expected output: 0 5 1 6 2 3 4
  for (int i = 0; i < v.size(); i++) cout << v[i] << " ";</pre>
  cout << endl;</pre>
  cout << suffix.LongestCommonPrefix(0, 2) << endl;</pre>
// BEGIN CUT
#endif
// END CUT
```

5.2 Binary Indexed Tree

```
#include <iostream>
using namespace std;
#define LOGSZ 17
int tree[(1<<LOGSZ)+1];</pre>
int N = (1 << LOGSZ);
// add v to value at x
void set(int x, int v) {
  while (x \le N)
   tree[x] += v;
   x += (x & -x);
// get cumulative sum up to and including x
int get(int x) {
 int res = 0;
  while(x) {
   res += tree[x];
   x -= (x & -x);
  return res:
// get largest value with cumulative sum less than or equal to x;
// for smallest, pass x-1 and add 1 to result
int getind(int x) {
```

```
int idx = 0, mask = N;
while mask && idx < N) {
  int t = idx + mask;
  if(x >= tree[t]) {
    idx = t;
    x -= tree[t];
  }
  mask >>= 1;
}
return idx;
```

5.3 Union-find set

```
#include <bits/stdc++.h>
using namespace std;
struct UnionFind {
    vector(int> C;
    UnionFind(int n) : C(n) { for (int i = 0; i < n; i++) C[i] = i; }
    int find(int x) { return (C[x] = x) ? x : C[x] = find(C[x]); }
    void merge(int x, int y) { C[find(x)] = find(y); }
}</pre>
```

5.4 KD-tree

```
// A straightforward, but probably sub-optimal KD-tree implmentation
// that's probably good enough for most things (current it's a
// - constructs from n points in O(n 1g^2 n) time
// - handles nearest-neighbor query in O(lg n) if points are well
     distributed
// - worst case for nearest-neighbor may be linear in pathological
     case
// Sonny Chan, Stanford University, April 2009
#include <iostream>
#include <vector>
#include <limits>
#include <cstdlib>
using namespace std;
// number type for coordinates, and its maximum value
typedef long long ntype;
const ntype sentry = numeric_limits<ntype>::max();
// point structure for 2D-tree, can be extended to 3D
struct point {
    ntype x, y;
    point(ntype xx = 0, ntype yy = 0) : x(xx), y(yy) {}
bool operator==(const point &a, const point &b)
    return a.x == b.x && a.y == b.y;
// sorts points on x-coordinate
bool on_x(const point &a, const point &b)
    return a.x < b.x;
// sorts points on y-coordinate
bool on_y(const point &a, const point &b)
    return a.y < b.y;</pre>
// squared distance between points
ntype pdist2(const point &a, const point &b)
    ntype dx = a.x-b.x, dy = a.y-b.y;
    return dx*dx + dy*dy;
// bounding box for a set of points
struct bbox
```

```
ntype x0, x1, y0, y1;
    bbox() : x0(sentry), x1(-sentry), y0(sentry), y1(-sentry) {}
    // computes bounding box from a bunch of points
    void compute(const vector<point> &v) {
        for (int i = 0; i < v.size(); ++i) {</pre>
            x0 = min(x0, v[i].x); x1 = max(x1, v[i].x);
            y0 = min(y0, v[i].y); y1 = max(y1, v[i].y);
    // squared distance between a point and this bbox, 0 if inside
    ntype distance(const point &p) {
        if (p.x < x0) {
    if (p.y < y0)
                                 return pdist2(point(x0, y0), p);
            else if (p.y > y1) return pdist2(point(x0, y1), p);
                                 return pdist2(point(x0, p.y), p);
        else if (p.x > x1) {
                                 return pdist2(point(x1, y0), p);
            if (p.y < y0)
            else if (p.y > y1) return pdist2(point(x1, y1), p);
            else
                                 return pdist2(point(x1, p.y), p);
        else
            if (p.y < y0)
                                 return pdist2(point(p.x, y0), p);
            else if (p.y > y1) return pdist2(point(p.x, y1), p);
                                 return 0:
            else
};
// stores a single node of the kd-tree, either internal or leaf
    bool leaf;
                    // true if this is a leaf node (has one point)
    point pt;
                    // the single point of this is a leaf
    bbox bound;
                    // bounding box for set of points in children
    kdnode *first, *second; // two children of this kd-node
    kdnode() \;\;:\;\; leaf(\textbf{false})\,,\;\; first(0)\,,\;\; second(0) \;\; \{\}
    "kdnode() { if (first) delete first; if (second) delete second; }
    // intersect a point with this node (returns squared distance)
    ntype intersect(const point &p) {
        return bound.distance(p);
    // recursively builds a kd-tree from a given cloud of points
    void construct(vector<point> &vp)
        // compute bounding box for points at this node
        bound.compute(vp);
        // if we're down to one point, then we're a leaf node
        if (vp.size() == 1) {
            leaf = true;
            pt = vp[0];
        else {
             // split on x if the bbox is wider than high (not best heuristic...)
            if (bound.x1-bound.x0 >= bound.y1-bound.y0)
                sort(vp.begin(), vp.end(), on_x);
            // otherwise split on y-coordinate
            else
                sort(vp.begin(), vp.end(), on_y);
            \ensuremath{//}\xspace divide by taking half the array for each child
            // (not best performance if many duplicates in the middle)
            int half = vp.size()/2;
            vector<point> vl(vp.begin(), vp.begin()+half);
            vector<point> vr(vp.begin()+half, vp.end());
            first = new kdnode(); first->construct(v1);
            second = new kdnode(); second->construct(vr);
// simple kd-tree class to hold the tree and handle queries
struct kdtree
    kdnode *root:
    // constructs a kd-tree from a points (copied here, as it sorts them)
    kdtree(const vector<point> &vp) {
        vector<point> v(vp.begin(), vp.end());
        root = new kdnode();
        root->construct(v);
    "kdtree() { delete root; }
```

```
// recursive search method returns squared distance to nearest point
    ntype search(kdnode *node, const point &p)
            // commented special case tells a point not to find itself
             if (p == node->pt) return sentry;
               return pdist2(p, node->pt);
       ntype bfirst = node->first->intersect(p);
       ntype bsecond = node->second->intersect(p);
        \ensuremath{//} choose the side with the closest bounding box to search first
        // (note that the other side is also searched if needed)
       if (bfirst < bsecond) {
            ntype best = search(node->first, p);
           if (bsecond < best)</pre>
               best = min(best, search(node->second, p));
           return best;
       else {
            ntype best = search(node->second, p);
           if (bfirst < best)</pre>
               best = min(best, search(node->first, p));
           return best;
    // squared distance to the nearest
   ntype nearest(const point &p) {
       return search (root, p);
};
// some basic test code here
int main()
    // generate some random points for a kd-tree
   vector<point> vp;
for (int i = 0; i < 100000; ++i) {</pre>
       vp.push_back(point(rand()%100000, rand()%100000));
    kdtree tree(vp);
    // query some points
    for (int i = 0; i < 10; ++i) {
       point q(rand()%100000, rand()%100000);
        return 0:
```

5.5 Splay tree

```
#include <cstdio>
#include <algorithm>
using namespace std;
const int N_MAX = 130010;
const int oo = 0x3f3f3f3f;
struct Node
  Node *ch[2], *pre;
  int val, size;
  bool isTurned;
} nodePool[N_MAX], *null, *root;
Node *allocNode(int val)
  static int freePos = 0;
  Node *x = &nodePool[freePos ++];
  x->val = val, x->isTurned = false;
  x->ch[0] = x->ch[1] = x->pre = null;
  x->size = 1;
  return x;
inline void update (Node *x)
  x->size = x->ch[0]->size + x->ch[1]->size + 1;
```

```
inline void makeTurned(Node *x)
  if(x == null)
    return;
  swap(x->ch[0], x->ch[1]);
x->isTurned ^= 1;
inline void pushDown(Node *x)
  if(x->isTurned)
    makeTurned(x->ch[0]):
    makeTurned(x->ch[1]);
    x->isTurned ^= 1;
inline void rotate(Node *x, int c)
  Node *y = x->pre;
  x->pre = y->pre;
  if(y->pre != null)
    y->pre->ch[y == y->pre->ch[1]] = x;
  y->ch[!c] = x->ch[c];
if(x->ch[c] != null)
  x->ch[c]->pre = y;
x->ch[c] = y, y->pre = x;
  update(y);
  if(y == root)
    root = x;
void splay(Node *x, Node *p)
  while(x->pre != p)
    if(x->pre->pre == p)
      rotate(x, x == x->pre->ch[0]);
    else
      Node *y = x->pre, *z = y->pre;
if(y == z->ch[0])
        if(x == y->ch[0])
           rotate(y, 1), rotate(x, 1);
        else
           rotate(x, 0), rotate(x, 1);
      else
        if(x == y->ch[1])
           rotate(y, 0), rotate(x, 0);
        else
          rotate(x, 1), rotate(x, 0);
  update(x);
void select(int k, Node *fa)
  Node *now = root;
  while (1)
    pushDown (now);
    int tmp = now->ch[0]->size + 1;
if(tmp == k)
      break;
    else if(tmp < k)</pre>
      now = now -> ch[1], k -= tmp;
    else
      now = now -> ch[0];
  splay(now, fa);
Node *makeTree(Node *p, int 1, int r)
  if(1 > r)
   return null;
  int \ mid = (1 + r) / 2;
  Node *x = allocNode(mid);
  x->pre = p:
  x\rightarrow ch[0] = makeTree(x, 1, mid - 1);
  x\rightarrow ch[1] = makeTree(x, mid + 1, r);
  update(x);
  return x;
```

```
int main()
  int n, m;
 null = allocNode(0);
  null->size = 0;
  root = allocNode(0);
  root->ch[1] = allocNode(oo);
  root->ch[1]->pre = root;
  update(root);
  scanf("%d%d", &n, &m);
  root->ch[1]->ch[0] = makeTree(root->ch[1], 1, n);
  splay(root->ch[1]->ch[0], null);
  while (m --)
   int a, b;
   scanf("%d%d", &a, &b);
    a ++, b ++;
    select(a - 1, null);
    select(b + 1, root);
    makeTurned(root->ch[1]->ch[0]);
  for(int i = 1; i <= n; i ++)</pre>
   select(i + 1, null);
printf("%d ", root->val);
```

5.6 Lowest common ancestor

```
const int max_nodes, log_max_nodes;
int num_nodes, log_num_nodes, root;
vector<int> children[max_nodes];
                                        // children[i] contains the children of node i
                                        // A[i][j] is the 2^j-th ancestor of node i, or -1 if that
int A[max_nodes][log_max_nodes+1];
      ancestor does not exist
int L[max nodes]:
                                        // L[i] is the distance between node i and the root
// floor of the binary logarithm of n
int lb(unsigned int n)
        return -1;
    int p = 0;
    if (n >= 1<<16) { n >>= 16; p += 16;
    if (n >= 1<< 8) { n >>= 8; p += 8;
    if (n >= 1 << 4) { n >>= 4; p += 4;
    if (n >= 1<< 2) { n >>= 2; p += 2;
   if (n >= 1<< 1) {
    return p;
void DFS(int i, int 1)
    L[i] = 1;
    for(int j = 0; j < children[i].size(); j++)</pre>
        DFS(children[i][j], 1+1);
int LCA(int p, int q)
     // ensure node p is at least as deep as node q
   if(L[p] < L[q])
        swap(p, q);
    // "binary search" for the ancestor of node p situated on the same level as {f q}
    for (int i = log_num_nodes; i >= 0; i--)
        if(L[p] - (1<<i) >= L[q])
           p = A[p][i];
    if(p == q)
        return p;
     // "binary search" for the LCA
    for(int i = log_num_nodes; i >= 0; i--)
        if(A[p][i] != -1 && A[p][i] != A[q][i])
            p = A[p][i];
            q = A[q][i];
    return A[p][0];
```

```
int main(int argc,char* argv[])
     // read num_nodes, the total number of nodes
    log_num_nodes=1b(num_nodes);
    for(int i = 0; i < num_nodes; i++)</pre>
        // read p, the parent of node i or -1 if node i is the root
        if(p != -1)
             children[p].push_back(i);
        else
             root = i;
     // precompute A using dynamic programming
    for(int j = 1; j <= log_num_nodes; j++)
    for(int i = 0; i < num_nodes; i++)</pre>
             if(A[i][j-1] != -1)
                 A[i][j] = A[A[i][j-1]][j-1];
             else
                 A[i][j] = -1;
    // precompute L
    DFS(root, 0);
    return 0:
```

6 Miscellaneous

6.1 Longest increasing subsequence

```
// Given a list of numbers of length n, this routine extracts a
// longest increasing subsequence.
// Running time: O(n log n)
    INPUT: a vector of integers
    OUTPUT: a vector containing the longest increasing subsequence
#include <iostream>
#include <vector>
#include <algorithm>
using namespace std;
typedef vector<int> VI;
typedef pair<int, int> PII;
typedef vector<PII> VPII;
#define STRICTLY_INCREASNG
VI LongestIncreasingSubsequence(VI v) {
  VPII best;
 VI dad(v.size(), -1);
  for (int i = 0; i < v.size(); i++) {</pre>
#ifdef STRICTLY_INCREASNG
    PII item = make_pair(v[i], 0);
    VPII::iterator it = lower_bound(best.begin(), best.end(), item);
    item.second = i;
    PII item = make_pair(v[i], i);
    VPII::iterator it = upper_bound(best.begin(), best.end(), item);
#endif
    if (it == best.end()) {
      dad[i] = (best.size() == 0 ? -1 : best.back().second);
      best.push_back(item);
    } else {
      dad[i] = it == best.begin() ? -1 : prev(it)->second;
      *it = item:
  VI ret;
  for (int i = best.back().second; i >= 0; i = dad[i])
   ret.push_back(v[i]);
  reverse(ret.begin(), ret.end());
  return ret;
```

6.2 Dates

```
// Routines for performing computations on dates. In these routines,
// months are expressed as integers from 1 to 12, days are expressed
// as integers from 1 to 31, and years are expressed as 4-digit
// integers.
#include <iostream>
#include <string>
using namespace std;
string dayOfWeek[] = {"Mon", "Tue", "Wed", "Thu", "Fri", "Sat", "Sun"};
// converts Gregorian date to integer (Julian day number)
int dateToInt (int m, int d, int y) {
 return
    1461 * (y + 4800 + (m - 14) / 12) / 4 +
    367 * (m - 2 - (m - 14) / 12 * 12) / 12 -
    3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +
    d - 32075;
// converts integer (Julian day number) to Gregorian date: month/day/year
void intToDate (int jd, int &m, int &d, int &y) {
  x = jd + 68569;
 n = 4 * x / 146097;

x -= (146097 * n + 3) / 4;

i = (4000 * (x + 1)) / 1461001;

x -= 1461 * i / 4 - 31;
  j = 80 * x / 2447;
  d = x - 2447 * j / 80;
 x = j / 11;

m = j + 2 - 12 * x;
  y = 100 * (n - 49) + i + x;
// converts integer (Julian day number) to day of week
string intToDay (int jd) {
  return dayOfWeek[jd % 7];
int main (int argc, char **argv) {
  int jd = dateToInt (3, 24, 2004);
  int m, d, y;
  intToDate (jd, m, d, y);
  string day = intToDay (jd);
       2453089
       3/24/2004
      Wed
  cout << jd << endl
    << m << "/" << d << "/" << y << endl
    << day << endl;
```

6.3 Prime numbers

```
// O(sqrt(x)) Exhaustive Primality Test
#include <cmath>
#define EPS 1e-7
typedef long long LL:
bool IsPrimeSlow (LL x)
  if(x<=1) return false;</pre>
  if(x<=3) return true;</pre>
  if (!(x%2) || !(x%3)) return false;
  LL s=(LL) (sqrt((double)(x))+EPS);
  for (LL i=5; i <= s; i+=6)
    if (!(x%i) || !(x%(i+2))) return false;
  return true;
// Primes less than 1000:
                               11
59
                                      13
       41
             43
                         53
                                     61
                                            67
                                                        73
                                                              79
                                                                     83
                                                                           89
                        107
                              109
                                    113
                                                                    149
                 103
                                                             139
                                                                          151
                 167
                        173
                              179
                                    181
251
                                          191
                                                 193
                                                       197
                                                             199
                                                                   211
277
            163
                  233
                        239
                              241
                                           257
                                                 263
                                                       269
                                                             271
            229
```

```
389
      439
            443
                   449
                         457
                                461
                                      463
                                             467
                                                                491
                         541
                                547
                                      557
                                             563
                                                   569
      599
                                      619
                                             631
                                                   641
                                                          643
                                                                647
                          683
                                691
                                      701
                                             709
                                                   719
                                                                733
                                                                       739
                                                                             743
      751
            757
                   761
                         769
                                773
                                      787
                                             797
                                                   809
                                                          811
                                                                821
                                                                             827
      829
            839
                   853
                         857
                                859
                                      863
                                            877
                                                   881
                                                          883
                                                                887
                                                                      907
                                                                             911
                         941 947
                                      953
// Other primes:
      The largest prime smaller than 10 is 7.
      The largest prime smaller than 100 is 97.
      The largest prime smaller than 1000 is 997.
      The largest prime smaller than 10000 is 9973.
      The largest prime smaller than 100000 is 99991.
The largest prime smaller than 1000000 is 999983.
      The largest prime smaller than 10000000 is 9999991.
      The largest prime smaller than 100000000 is 99999989.
      The largest prime smaller than 1000000000 is 999999937.
      The largest prime smaller than 10000000000 is 9999999967.
      The largest prime smaller than 10000000000 is 99999999977.
      The largest prime smaller than 100000000000 is 99999999989. The largest prime smaller than 100000000000 is 999999999971.
      The largest prime smaller than 10000000000000 is 9999999999973. The largest prime smaller than 10000000000000 is 999999999999989.
      The largest prime smaller than 100000000000000 is 99999999999937.
      The largest prime smaller than 1000000000000000 is 999999999999997.
```

6.4 Miller-Rabin Primality Test

```
// Randomized Primality Test (Miller-Rabin):
    Error rate: 2^(-TRIAL)
// Almost constant time. srand is needed
#include <stdlib h>
#define EPS 1e-7
typedef long long LL:
LL ModularMultiplication (LL a, LL b, LL m)
        LL ret=0, c=a;
        while(b)
                if(b&1) ret=(ret+c)%m;
        return ret;
LL ModularExponentiation(LL a, LL n, LL m)
        LL ret=1, c=a;
        while(n)
                if(n&1) ret=ModularMultiplication(ret, c, m);
                n>>=1; c=ModularMultiplication(c, c, m);
        return ret;
bool Witness(LL a, LL n)
  int t=0;
        while(!(u&1)){u>>=1; t++;}
        LL x0=ModularExponentiation(a, u, n), x1;
        for(int i=1;i<=t;i++)
                x1=ModularMultiplication(x0, x0, n);
                if (x1==1 && x0!=1 && x0!=n-1) return true;
                x0=x1:
        if(x0!=1) return true;
        return false;
LL Random(LL n)
  LL ret=rand(); ret *= 32768;
        ret+=rand(); ret*=32768;
        ret+=rand(); ret*=32768;
        ret+=rand();
  return ret%n:
bool IsPrimeFast (LL n, int TRIAL)
  while (TRIAL--)
```

```
LL a=Random(n-2)+1;
  if(Witness(a, n)) return false;
}
return true;
```

6.5 Fast exponentiation

```
Uses powers of two to exponentiate numbers and matrices. Calculates
n^k in O(log(k)) time when n is a number. If A is an n x n matrix.
calculates A^k in O(n^3*log(k)) time.
#include <iostream>
#include <vector>
using namespace std;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
T power(T x, int k) {
  T ret = 1;
  while(k) {
    if(k & 1) ret *= x;
    k >>= 1; x *= x;
  return ret;
VVT multiply(VVT& A, VVT& B) {
  int n = A.size(), m = A[0].size(), k = B[0].size();
  VVT C(n, VT(k, 0));
  for (int i = 0; i < n; i++)
   for(int j = 0; j < k; j++)
for(int l = 0; l < m; l++)
        C[i][j] += A[i][1] * B[1][j];
  return C:
VVT power(VVT& A, int k) {
  int n = A.size();
  VVT ret(n, VT(n)), B = A;
  for(int i = 0; i < n; i++) ret[i][i]=1;</pre>
  while(k) {
    if(k & 1) ret = multiply(ret, B);
    k >>= 1; B = multiply(B, B);
  return ret:
int main()
  /* Expected Output:
     2.37^48 = 9.72569e+17
     376 264 285 220 265
     550 376 529 285 484
     484 265 376 264 285
285 220 265 156 264
     529 285 484 265 376 */
  double n = 2.37:
  int k = 48;
  cout << n << "^" << k << " = " << power(n, k) << endl;
  double At [5] [5] = {
    { 0, 0, 1, 0, 0 },
     { 1, 0, 0, 1, 0 },
    { 0, 0, 0, 0, 1 },
    { 1, 0, 0, 0, 0 },
    { 0, 1, 0, 0, 0 } };
  vector <vector <double> > A(5, vector <double>(5));
  for(int i = 0; i < 5; i++)
for(int j = 0; j < 5; j++)
A[i][j] = At[i][j];</pre>
  vector <vector <double> > Ap = power(A, k);
  cout << endl;
```

```
for(int i = 0; i < 5; i++) {
  for(int j = 0; j < 5; j++)
    cout << Ap[i][j] << " ";
  cout << endl;
}</pre>
```

6.6 C++ input/output

```
#include <bits/stdc++.h>
using namespace std:
int main()
    // Ouput a specific number of digits past the decimal point,
    // in this case 5
    cout.setf(ios::fixed); cout << setprecision(5);</pre>
    cout << 100.0/7.0 << endl;</pre>
    cout.unsetf(ios::fixed);
    // Output the decimal point and trailing zeros \,
    cout.setf(ios::showpoint);
    cout << 100.0 << end1:
    cout.unsetf(ios::showpoint);
    // Output a '+' before positive values
    cout.setf(ios::showpos);
    cout << 100 << " " << -100 << endl;
    cout.unsetf(ios::showpos);
    // Output numerical values in hexadecimal
    cout << hex << 100 << " " << 1000 << " " << 10000 << dec << endl;
```

6.7 Knuth-Morris-Pratt

```
Finds all occurrences of the pattern string p within the
text string t. Running time is O(n + m), where n and m
are the lengths of p and t, respecitvely.
#include <iostream>
#include <string>
#include <vector>
using namespace std;
typedef vector<int> VI;
void buildPi(string& p, VI& pi)
  pi = VI(p.length());
  int k = -2;
  for(int i = 0; i < p.length(); i++) {</pre>
    while (k \ge -1 \&\& p[k+1] != p[i])
      k = (k == -1) ? -2 : pi[k];
    pi[i] = ++k;
int KMP (string& t. string& p)
  VI pi;
 buildPi(p, pi);
  int k = -1;
  for(int i = 0; i < t.length(); i++) {</pre>
    while (k \ge -1 \&\& p[k+1] != t[i])
     k = (k == -1) ? -2 : pi[k];
    if(k == p.length() - 1) {
     // p matches t[i-m+1, ..., i]
      cout << "matched at index " << i-k << ": ";
      cout << t.substr(i-k, p.length()) << endl;</pre>
     k = (k == -1) ? -2 : pi[k];
  return 0:
int main()
```

```
string a = "AABAACAADAABAABA", b = "AABA";
KMP(a, b); // expected matches at: 0, 9, 12
return 0;
```

6.8 Latitude/longitude

```
Converts from rectangular coordinates to latitude/longitude and vice
versa. Uses degrees (not radians).
#include <iostream>
#include <cmath>
using namespace std;
struct 11
  double r, lat, lon;
};
struct rect
  double x, y, z;
};
11 convert(rect& P)
  11 Q;
  Q.r = sqrt(P.x*P.x+P.y*P.y+P.z*P.z);
Q.lat = 180/M_PI*asin(P.z/Q.r);
  Q.lon = 180/M_PI*acos(P.x/sqrt(P.x*P.x+P.y*P.y));
  return 0;
rect convert(ll& Q)
  P.x = Q.r*cos(Q.lon*M_PI/180)*cos(Q.lat*M_PI/180);
  P.y = Q.r*sin(Q.lon*M_PI/180)*cos(Q.lat*M_PI/180);
  P.z = Q.r*sin(Q.lat*M_PI/180);
  return P:
int main()
  rect A:
  11 B;
  A.x = -1.0; A.y = 2.0; A.z = -3.0;
  B = convert(A);
cout << B.r << " " << B.lat << " " << B.lon << endl;</pre>
  A = convert(B);
cout << A.x << " " << A.y << " " << A.z << endl;
```

6.9 Vim settings

```
set enc=utf-8
set fenc=utf-8
scriptencoding utf-8
```

```
set fencs=utf-8,ucs-bom,gb18030,gbk,gb2312,cp936
syntax on
filetype plugin on
filetype indent on
set mouse=a
set nocompatible
set tabstop=4
set shiftwidth=4
set expandtab
set smarttab
set autoindent
set textwidth=1000
set showmatch
set ruler
set hlsearch
set incsearch
set ignorecase
set smartcase
set number
set relativenumber
set fdm=marker
set scrolloff=10
set showcmd
set backspace=indent,eol,start
set formatoptions=q,r
set nowrap
set foldmethod=indent
set foldlevelstart=10
set wildmenu
set omnifunc=syntaxcomplete#Complete
"the trail config failed in generate latex "remember to add in the site"
set list listchars=tab:>-,trail:
set t_Co=256
set background=dark
nnoremap ; :
inoremap <silent> jj <ESC>
nnoremap <silent> <HOME> 1
inoremap <silent> <HOME> <ESC>^i
vnoremap <silent> <HOME> 0w
nnoremap <silent> <F9> :set nowrap<CR>
nnoremap <silent> <F10> :set wrap<CR>
```