

# Software Specification

---

## Dijkstra's Algorithm Implementation

=====

### Introduction

-----

Dijkstra's algorithm is a well-known graph search algorithm used to find the shortest

### Implementation

-----

### Graph Representati

will use an adjacency list to represent the graph, where each node is associated w

### Node and Edge Manageme

will use the NetworkX library to add, remove, and update nodes and edges in the g

### Shortest Path Calculati

will use Dijkstra's algorithm to calculate the shortest path between two node

### Graph Visualizati

# Software Specification

---

will use Matplotlib to visualize the shortest path between two node

## Code

--

`pyth

port networkx as

port matplotlib.pyplot as p

port hea

ass Grap

f \_\_init\_\_(self

If.G = nx.DiGraph

f add\_node(self, node\_id

If.G.add\_node(node\_i

f add\_edge(self, node1, node2, weight

If.G.add\_edge(node1, node2, weight=weigh

f dijkstra(self, start\_node, end\_node

## Software Specification

---

initialize distances and previous node

```
distances = {node: float('inf') for node in self.G.nodes()}
```

```
previous_nodes = {node: None for node in self.G.nodes()}
```

```
distances[start_node] =
```

priority queue

```
priority_queue = [(0, start_node)]
```

while priority\_queue

```
current_distance, current_node = heapq.heappop(priority_queue)
```

if current distance is greater than already known distance

```
current_distance > distances[current_node]
```

return

iterate over neighbors

```
for neighbor in self.G.neighbors(current_node):
```

```
weight = self.G[current_node][neighbor]['weight']
```

```
distance = current_distance + weight
```

## Software Specification

---

date distance and previous node if shorter path is found

distance < distances[neighbor]

distances[neighbor] = distance

previous\_nodes[neighbor] = current\_node

heapq.heappush(priority\_queue, (distance, neighbor))

while shortest\_path is not None:

current\_node =

current\_node = end\_node

while current\_node is not None:

path.append(current\_node)

current\_node = previous\_nodes[current\_node]

path.reverse()

return distances[end\_node], path

def visualize(self, path):

G = nx.spring\_layout(self.G)

G.draw(self.G, pos, with\_labels=True, node\_color='lightblue')

## Software Specification

---

```
.draw_networkx_edges(self.G, pos, edgelist=[(path[i], path[i+1]) for i in range(len(p
```

```
t.show
```

```
ample usa
```

```
aph = Graph
```

```
aph.add_node(
```

```
aph.add_node(
```

```
aph.add_node(
```

```
aph.add_edge(1, 2,
```

```
aph.add_edge(1, 3,
```

```
aph.add_edge(2, 3,
```

```
stance, path = graph.dijkstra(1,
```

```
int(f"Shortest distance: {distance}
```

```
int(f"Shortest path: {path}
```

```
aph.visualize(pat
```

### Algorithm

---

# Software Specification

---

-----

**Initialize distances and previous nodes for all nodes in the graph**

**Set the distance of the start node to 0 and add it to the priority queue**

**While the priority queue is not empty**

**Dequeue the node with the smallest distance**

**If the current distance is greater than the already known distance, skip this node**

**Iterate over the neighbors of the current node**

**Calculate the tentative distance to the neighbor**

**If the calculated distance is smaller than the already known distance, update the distance**

**Add the neighbor to the priority queue**

**Build the shortest path by backtracking from the end node to the start node**

## Time Complexity

-----

**The time complexity of Dijkstra's algorithm is  $O(|E|\log|V|)$  in the worst case, where  $|E|$  is the number of edges and  $|V|$  is the number of vertices.**

## Space Complexity

-----

# Software Specification

---

The space complexity of Dijkstra's algorithm is  $O(|V| + |E|)$ , where  $|V|$  is the number of vertices and  $|E|$  is the number of edges.

Note: This implementation assumes that the graph is represented as an adjacency list.

*Generated on 2025-04-06 at 12:41:43*