DEPARTMENT OF ELECTRICAL AND ELECTRONIC ENGINEERING **EXAMINATIONS 2011** 

EEE/ISE PART II: MEng, BEng and ACGI

## **COMMUNICATIONS 2**

Wednesday, 15 June 2:00 pm

Time allowed: 2:00 hours

There are THREE questions on this paper.

Answer ALL questions. Q1 carries 40% of the marks. Questions 2 and 3 carry equal marks (30% each).

Any special instructions for invigilators and information for candidates are on page 1.

Examiners responsible

First Marker(s): C. Ling

Second Marker(s): J.A. Barria

## **EXAM QUESTIONS**

- 1. a) i) Explain the terms "noise", "external noise", and "internal noise". [3]
  - ii) Explain the terms "white noise", "Gaussian noise", and "additive white Gaussian noise". [3]
  - iii) Consider bandpass noise  $n(t) = n_c(t) \cos(2\pi f_c t) n_s(t) \sin(2\pi f_c t)$  which has the power spectral density shown in Fig. 1.1. Draw the power spectral density (PSD) of baseband noise  $n_c(t)$  or  $n_s(t)$  if the center frequency is chosen as:
    - $f_c = 7 \text{ Hz}$
    - $f_c = 10 \,\mathrm{Hz}$

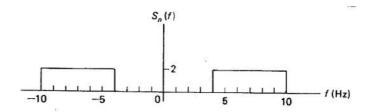


Figure 1.1 PSD of bandpass noise.

- b) i) Explain the advantages and disadvantages of synchronous detection and envelope detection respectively for standard amplitude modulation (AM).
  - ii) With help of a diagram, explain the operation of coherent detection for frequency shift keying (FSK). [5]
- c) i) Name two primary resources in communications. [2]
  - ii) Write down the expression of the capacity of a Gaussian channel with bandwidth W. [2]
  - iii) What does the channel coding theorem say about the relation between transmission rate *R* and channel capacity *C*? [3]
  - iv) Consider a telephone line channel. If the signal to noise ratio (SNR) is 20 dB and the bandwidth available is 4 kHz, calculate the corresponding channel capacity. [3]

d) Consider an information source generating the random variable *X* with probability distribution

				<i>x</i> <sub>4</sub>
$P(X=x_k)$	0.4	0.2	0.25	0.15

- i) Calculate the entropy of this source. [3]
- ii) Construct a binary Huffman code. [5]
- iii) Compute the average codeword length. [2]

2. a) Figure 2.1 shows the diagram of the FM receiver. The bandpass filter has bandwidth  $B_T$ , while the baseband low-pass filter has bandwith W. Let the FM signal be  $s(t) = A\cos[2\pi f_c t + 2\pi k_f \int_0^t m(\tau)d\tau]$  and assume the bandpass noise  $w(t) = n_c(t)\cos(2\pi f_c t) - n_s(t)\sin(2\pi f_c t)$  has single-sided power spectral density  $N_0$ .

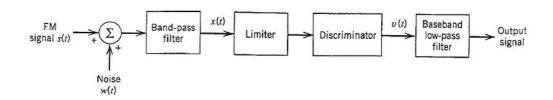


Figure 2.1 Block diagram of the FM receiver.

- i) Explain the function of each of the blocks in Figure 2.1. [5]
- ii) Given the noise at the output of the discriminator

$$f_i(t) = \frac{1}{2\pi A} \frac{dn_s(t)}{dt},$$

derive an expression for its power spectral density.

- [5]
- Sketch power spectral densities of  $n_s(t)$ ,  $f_i(t)$ , and the noise at the output of the lowpass filter. [5]
- b) Consider pre-emphasis and de-emphasis.
  - i) Show that in order to achieve flat noise power spectral density at the output of the FM receiver, the ideal de-emphasis filter has a transfer function  $H_{de}(f) = 1/f$  within the message bandwidth. [5]
  - ii) Discuss why an FM transmitter with a corresponding pre-emphasis filter  $H_{pre}(f) = f$  is essentially phase modulation. [5]
- c) The signal-to-noise ratio (SNR) improvement factor is defined as

$$I = \frac{\text{Noise power without pre/de-emphasis}}{\text{Noise power with pre/de-emphasis}}.$$

Derive an expression of the improvement factor I for the scaled ideal de-emphasis filter  $H_{de}(f) = f_0/f$ , then compute the gain in dB for the parameter W = 15 kHz and  $f_0 = 3 \text{ kHz}$ .

3. a) A uniform quantizer for PCM has  $2^n$  levels. The input signal is

$$m(t) = [A + A\cos(\omega_1 t)]\cos(\omega_2 t)$$

where  $\omega_1 \neq \pm \omega_2$ . Assume the dynamic range of the quantizer matches that of the input signal.

- i) Work out the signal power. [3]
- ii) Write down the probability density function of the quantization noise and the quantization noise power. [3]
- iii) Work out the SNR in dB at the output of the quantizer. [3]
- iv) Determine the minimum value of *n* such that the output SNR is no less than 60 dB. [3]
- v) What can be done to increase the output SNR? [3]
- b) An (n,k) linear block code has the following parity check matrix

$$H = \begin{bmatrix} 1 & 1 & 1 & 0 & 1 & 0 & 0 \\ 1 & 1 & 0 & 1 & 0 & 1 & 0 \\ 1 & 0 & 1 & 1 & 0 & 0 & 1 \end{bmatrix}$$
 (3.1)

- i) What are the values of n and k? [2]
- ii) Give a systematic generator matrix G of this code. [3]
- iii) Compute the syndrome table for a single error. [5]
- iv) The vector y = [1000001] is received. Find the syndrome and hence the most likely data bits. [5]

ANSWERS

B - Bookwork

E - New examples

A - New applications

EEZ-4 COMMUNICATIONS 2 (.a) i) Noise refers to unwanted waves that disturb Communications.

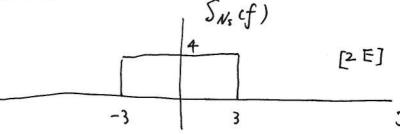
[3B]

External noise: interference from hearby channes, human-made noise, natural noise for external.

Internal noise: noise from within electronic devices, Such as thermal noise.

ii). White noise: the PSD is constant [3 B] Ganssian noise: the PDF is Ganssian. additive white Ganssian noise: noise is additive, white, and Gaussian.

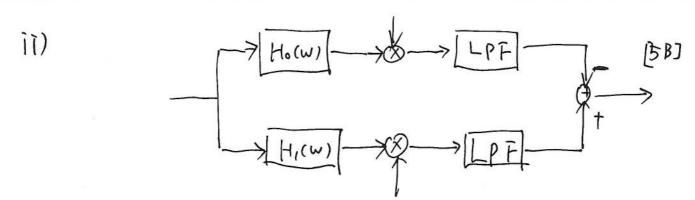
iii). fc= THz



fc = 10 Hz

b) i) Synchronous detection: good noise performance, complex sync circuit.

> Envelope detection: simple circuit, good performance at high SVR, but suffers from threshold effect. [3B]



The BPFs Ho(w) and H<sub>1</sub>(w) have central frequency found f<sub>1</sub>, respectively.

c) i) power, bandwidth

[28]

 $C = w \log_2(1 + SNR)$ 

[28]

iii) Channel coding theorem, As long as  $R \leq C$ , we can achieve reliable communication over a noisy channel (i.e., with arbitrarily small probability of error); conversely, it is impossible to transmit messages without error if R > C. [38]

iV) 20dB ⇒ 3NR=100

[3 E]

C = 4K × log\_2(1+100) = 26.6 kbps

d) i) 
$$H(x) = -\sum p(x_k) \log p(x_k)$$
 BE]  
=-0.4 × log 0.4 - 0.2 × log 0.2 - 0.25 × log 0.25 - 0.15 × log 0.15  
= 1.9

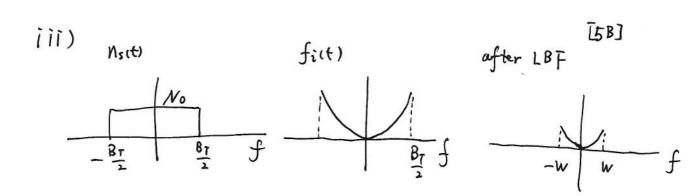
$$(iii)$$
 L =  $1 \times 0.4 + 2 \times 0.25 + 3 \times 0.35$  [2E]  
=  $1.95$ 

2. a) i) The bandpass filter removes out-of-band noise. The Limiter results in a constant envelope. The discriminator outputs the deviation in the instantaneous frequency, i.e., it recover the message signal.

Low-pass filter: has a bondwidth w. It passer the message and removes out-of-band noise.

ii) This can be viewed as a linear system with transfer function  $\frac{1}{2\pi A}$   $j_2\pi f = j\frac{f}{A}$ . [38]

Therefore, PSD for fitty is  $\frac{f^2}{A^2} N_0, \qquad |f| \leq \frac{Br}{2}.$ 



- b) i) It's clear that to equalize the PSD  $\frac{f^2}{A^2}N_0$ , we need a de-emphasis fitter  $H_{delf}$ ) =  $\frac{1}{f}$ . [5]
  - ii) Hpre(f) = f is a differentiation circuit. Thus the signal becomes  $Sct) = A cos[2\pi fet + kfmct)]$ , which is PM.

$$P_N = \int_{-w}^{w} \frac{f^2}{A^2} N_0 df = \frac{2}{3} \frac{w^3 N_0}{A^2}$$
 [18]

with deemphasis, the noise power is

$$P_{N} = \int_{-W}^{W} \frac{f^{2}}{A^{2}} N_{0} \cdot \frac{f^{2}}{f^{2}} df$$

$$= \int_{-W}^{W} \frac{f^{2}N_{0}}{A^{2}} df$$

$$= \frac{2W f^{2}N_{0}}{A^{2}} .$$
[2A]

Then,
$$\hat{I} = \frac{\frac{2 W^3 N_0}{3 A^2}}{\frac{2 W f_0^2 N_0}{A^2}} = \frac{W^2}{2 f_0^2}$$
[[A]

When w = 15 kHz, fo = 2.1 kHz

$$\bar{1} = \frac{15^2}{2 \cdot 2.1^2} = 25.5$$

In dB,

3. a) i) 
$$P_S = \frac{A^2}{2} + \frac{4^2}{4} = \frac{3}{4}A^2$$

[3 E]

[3 E]

ii) The quantization noise has a uniform PDF 
$$f(q) = \frac{1}{\Delta}, \quad |g| < \frac{\Delta}{2}$$

Power: 
$$P_N = \int_{-\frac{\alpha}{2}}^{\frac{\alpha}{2}} \chi^2 \cdot \frac{1}{\Delta} d\chi$$

$$= \frac{2}{\Delta} \cdot \frac{1}{3} \left(\frac{\Delta}{2}\right)^3$$

$$= \frac{\Delta^2}{12}$$

The signal range is 
$$[A,A]$$
.
$$\Delta = \frac{2A}{2^n} = \frac{A}{2^{n-1}}$$

$$SNR = \frac{P_S}{P_N} = \frac{\frac{3}{4}A^2}{\frac{(A/2^{n-1})^2}{12}} = \frac{9}{4} \cdot 2^{2n}$$
 [3E]

$$SNR(dB) = 6n + 3.5 dB$$

b) i) 
$$n=7$$
  $k=4$  [ZE]

ii)  $G = \begin{bmatrix} 1 & 0 & 0 & 1 & 1 & 1 \\ 0 & 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 & 0 & 1 \end{bmatrix}$ 

[3E]

$$S = y \cdot H^{T}$$

$$= (110)$$

$$\Rightarrow e = (0100000)$$

$$\Rightarrow \text{ the Second bit is wrong}$$

$$\Rightarrow \text{ sent code word } \chi = [1100001]$$

$$\Rightarrow \text{ data bits} = [1100]$$