

Sequences

Klaus Draeger

Lecture 7: Sequences

One property of sets is that all the following are equal:

$$\begin{aligned} & \{ 0, 1, 2, 5, 23, 99 \} \\ &= \{ 1, 99, 5, 2, 23, 0 \} \\ &= \{ 2, 1, 99, 5, 2, 23, 1, 0, 1 \} \end{aligned}$$

Sometimes it is necessary to:

- distinguish values of a set by **position**,
- or to **permit duplicate** values,
- or to impose some **ordering** on the values.

None of these can be achieved using a flat **set**.

In such cases a **sequence** is the correct structure to use.

Defining Sequences

- For any set **S** there is a set **seq(S)** of S-valued sequences.
- A sequence is defined by enclosing its values in **square brackets []**
- Example: for **Colours = {red, yellow, green}**
we have [red,green,red,red] : seq(Colours)
- Note that a sequence can be empty: []
- Like for functions, seq(S) has **specialised subsets**:
 - seq1(S) – the set of **non-empty** sequences
 - iseq(S) – the set of **injective** (non-repeating) sequences

Defining Sequences

- The defining property of sequences is that they are:
 - **Indexed** by numbers $1, \dots, n$ (for some n)
 - That is, they are **functions** $f : \mathbf{NAT1} \rightarrow \mathbf{S}$
with the additional constraint that **dom(f)** is an **initial segment** $1..n$ of $\mathbf{NAT1}$
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which we can simplify using **n = ???**

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- So we have the equality
$$\text{seq}(S) = \{f \mid f:\text{NAT1} \rightarrow S \ \& \ \exists(n). (n:\text{NAT} \ \& \ \text{dom}(f)=1..n)\}$$
which we can simplify using $n = \text{card}(f)$:
$$\text{seq}(S) = \{f \mid f:\text{NAT1} \rightarrow S \ \& \ \text{dom}(f)=1..\text{card}(f)\}$$

Working with Sequences: Access

- Since sequences are functions, we can get the **n**th element by **evaluating** at **n**
- If `colSequence = [red,green,green,blue,red]` then `colSequence(1) = red` and `colSequence(4) = blue`
- The **first** and **last** operators return the first/last element: `first(colSequence) = red`, `last(colSequence) = red`
- How do these differ from **min** and **max**?

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- The **first** and **last** operators return the first/last element: `first(colSequence) = red`, `last(colSequence) = red`
- In general, for any sequence **seq**, `first(seq) = seq(1)` and `last(seq) = seq(card(seq))`
- None of these are defined for the empty sequence `[]`.

Working with Sequences: Truncation

- **front** and **tail** complement the first and last operators:
 - **front(seq)** is **seq** with its last element removed
 - **tail(seq)** is **seq** with its first element removed
 - Examples:
 - **front([2,3,5,7,11]) = [2,3,5,7]**
 - **tail([2,3,5,7,11]) = [3,5,7,11]**
 - Like first and last, they are undefined for the empty sequence.
- How could we re-create these if they did not exist?
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- How could we re-create these if they did not exist?
 - To define front(seq): **restrict** seq removing the last index
 $\text{front}(\text{seq}) = (\text{dom}(\text{seq}) - \text{card}(\text{seq})) <| \text{seq}$
e.g. $\text{front}([2,3,5,7,11]) = (1..4) <| [2,3,5,7,11]$
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 - To define **tail(seq)**: apply a **shift** (a restriction of **succ**)
 $\text{tail}(\text{seq}) = (\text{dom}(\text{seq}) <| \text{succ}) ; \text{seq}$
e.g. $\text{tail}([2,3,5,7,11]) = [2,3,4,5] ; [2,3,5,7,11]$

Working with Sequences: Insertion

- The opposite of the truncations are the front and end **insertions**.
- $\mathbf{x \rightarrow seq}$ is the sequence **seq** with **x** inserted at the **front**
so \rightarrow is a function in $(\mathbf{X * seq(X)}) \rightarrow \mathbf{seq(X)}$
with $\mathbf{x \rightarrow [z_1, \dots, z_n] = [x, z_1, \dots, z_n]}$
- $\mathbf{seq <- y}$ is the sequence **seq** with **y** inserted at the **back**
so \leftarrow is a function in $(\mathbf{seq(X) * X}) \rightarrow \mathbf{seq(X)}$
with $\mathbf{[z_1, \dots, z_n] <- y = [z_1, \dots, z_n, y]}$
- This lets us express the relationship between first/last/front/tail as
 - $\mathbf{seq = front(seq) <- last(seq)}$
 - $\mathbf{seq = first(seq) \rightarrow tail(seq)}$

Sequence Example: Stack

- Our goal is to model a **stack** of natural numbers.
- It should have the following operations:
 - **Push** – pushes a number onto the stack
 - **Pop** – pops (i.e. removes) and returns the number at the top of the stack
 - this requires the stack to be **non-empty**
 - **IsEmpty** – returns **Yes** if the stack is empty; otherwise returns **No**
- To do this, we need to
 - represent the stack itself
 - a sequence works for this
 - Figure out how to perform the operations

Sequence Example: Stack

MACHINE Stack

SETS

ANSWER = { Yes, No }

VARIABLES

stack

INVARIANT

stack : seq(NAT)

INITIALISATION

stack := []

OPERATIONS

res <-- isEmpty = ?

Sequence Example: Stack

```
//...
```

```
res <-- isEmpty =  
  IF stack = []  
  THEN  
    res := Yes  
  ELSE  
    res := No  
  END  
END;
```

```
push(num) = ?
```

```
pop = ?
```

Sequence Example: Stack

```
// using the back of the sequence as the top
```

```
push(num) =  
  PRE  
    num : NAT  
  THEN  
    stack := stack <- num  
  END;  
  
res <-- pop =  
  PRE  
    stack /= []  
  THEN  
    stack := front(stack) ||  
    res := last(stack)  
  END  
END
```


Doing More with Sequences

- Beyond the operators discussed so far, we may want to:
 - **Combine** two sequences
For example, combine [1,2,3] and [3,2,1] into [1,2,3,3,2,1]
Similar to **concatenating** two strings in Java
 - **Split** a sequence at some position
For example, split [1,2,3,4,5] into [1,2,3] and [4,5]
 - **Reverse** a sequence

These are built-in, and we will see how to define others:

- **Insert** or **delete** values in anywhere
- **Filter** a sequence

- The concatenation operator \wedge combines two sequences:
 $[1,3,5] \wedge [2,4,5] = [1,3,5,2,4,5]$
- Note how the insertions are special cases:
 - $x \rightarrow \text{seq} = [x] \wedge \text{seq}$
 - $\text{seq} \leftarrow z = \text{seq} \wedge [z]$
- Why not just use the existing union ($\setminus/\$) operator?

Concatenation

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 - $x \rightarrow \text{seq} = [x] \wedge \text{seq}$
 - $\text{seq} \leftarrow z = \text{seq} \wedge [z]$
- Why not just use the existing union (\setminus/\setminus) operator?
 - \setminus/\setminus works at the set level
 - At this level $[1,3,5]$ and $[2,4,5]$ are simply $\{1|->1,2|->3,3|->5\}$ and $\{1|->2,2|->4,3|->5\}$
 - Their union is $\{1|->1,2|->3,3|->5,1|->2,2|->4,3|->5\}$
This is not a function (e.g. $1|->1$ **and** $1|->2$)
and therefore also not a sequence

We **can** re-define \wedge though, see next tutorial!

Splitting sequences

- The operators $/\backslash$ (up-arrow) and $\backslash/$ (down-arrow) split a sequence at a given position:
For any n between 0 and $\text{card}(\text{seq})$,
 - $\text{seq} / \backslash n$ returns the first n elements of seq
 - $\text{seq} \backslash / n$ returns the rest of seq
i.e. everything except the first n elements
- If $c = [\text{red}, \text{orange}, \text{yellow}, \text{green}, \text{blue}, \text{purple}]$:
 - $c / \backslash 2 = [\text{red}, \text{orange}]$
 - $c \backslash / 2 = [\text{yellow}, \text{green}, \text{blue}, \text{purple}]$
- So for all n , $(\text{seq} / \backslash n) \wedge (\text{seq} \backslash / n) = \text{seq}$
- Again, let's see how we could re-define these ourselves:

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- So for all n , $(\text{seq} / \backslash n) \wedge (\text{seq} \backslash / n) = \text{seq}$
- Again, let's see how we could re-define these ourselves:
 - $\text{seq} / \backslash n = (1..n) <| \text{seq}$
 - $\text{seq} \backslash / n$: see tutorial!

Adding and Removing

- With these operators in place we can insert values wherever we want:
- To **insert** a value at position **n** in **seq**,
 - Take the first **n-1** elements of **seq**,
 - add the new value,
 - then append the remaining values of **seq**.

e.g. we can insert **3** in the middle of **seq = [1,2,4,5]** using
 $((seq//2)<-3) \wedge (seq\|/2) = ([1,2]<-3) \wedge [4,5] = [1,2,3,4,5]$

- To **remove** the element at position **n**,
 - Take the first **n-1** elements of **seq**,
 - and append the values of **seq** past the **nth**:

To remove the 4th entry in **seq = [1,2,3,4,5]**,
use $(seq//\backslash 3) \wedge (seq\backslash\|/4) = [1,2,3] \wedge [5] = [1,2,3,5]$

- Often we want to remove elements based on their **value**, not **position**
- For example, for a list of numbers **seq** = **[1,2,3,4,3,2,1,4,2,1,3]**
filter out the even numbers,
leaving just the odd ones: **[1,3,3,1,1,3]**
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- So suppose we have a sequence **seq** and subset **filter**, how to define a restriction of **seq** to values from **filter**?
 - Get the set of indices for elements in the filter
e.g. the indices of odd numbers are {1,3,5,7,10,11}
 - Define a sequence **select** containing them in order
e.g. [1,3,5,7,10,11]
 - Compose them (select ; seq)

We will do this in ProB.

Filtering

