

# ACM-ICPC World Finals 2019

Team Reference Document

University of Illinois at Urbana-Champaign VIM - Help poor children!

## Coach

Alan Mattox Beckman Jr

## Contestants

Lan Dao, Jiacheng Liu, Zexuan Zhong

$\mathbf{C}$	Contents		7	Ma	auth	14
				7.1		14
1	Getting Started	<b>2</b>				14
	1.1 Vimrc	2			1	14
	1.2 Starter Code	2				
	1.3 C++ Grammar, STL	2			I I I I I I I I I I I I I I I I I I I	15
					1	15
<b>2</b>	Data Structures	<b>2</b>				15
	2.1 Binary Indexed Tree 2D	2			7.1.7 Discrete log	16 16
	2.2 Segment Tree 2D	3			7.1.9 Euler function	
	2.3 Persistent Segment Tree	3				16
	2.4 Treap	4		7.2		17
	2.5 Splay Tree	4		,		17
	2.6 Mo's Algorithm	6				17
					7.2.3 Gaussian Elimination	17
3	Graph Theory	6		7.3		17
	3.1 Ford Fulkerson	6				17
	3.2 Dinic	6				
	3.3 Tarjan	7			7.3.3 Number Theoretic Transform	
	3.4 Topo Sort	8		7.4	FFT	19
	3.5 2-SAT	8	8	Geo	ometry	19
	3.6 Lowest Common Ancestor - $O(n \log n)$	8	Ü	8.1	•	19
	3.7 Eulerian Circuit	8		8.2		21
	3.8 Dominator Tree	9		8.3	Halfplane	22
	3.9 Centroid Decomposition	9		8.4	Polygon	23
	3.10 Heavy Light Decomposition	10		8.5	Circle	24
	5.10 Heavy Light Decomposition	10		8.6	Simplex volume	24
1	Dynamic Programming	10		8.7	Count gridpoints under a line	25
-	4.1 Convex Hull Trick	10		8.8	Simpson's Union Of Circles	25
	4.2 Dynamic Convex Hull Trick	11	9	Mis	S.C.	26
	4.2 Dynamic Convex Itun Itick	11	J		Date	26
5	String	11		0.1	9.1.1 Date to Day of Week	-
	5.1 Z-Function	11			9.1.2 Count Days from AD	
	5.2 Suffix Array				v	
	5.3 Aho-Corasick Automata					
	5.4 Palindromic Tree	13				
	o.i i dimendime iree	10				
6	Game Theory	14				
	6.1 Nim Product	14				

## 1 Getting Started

#### 1.1 Vimrc

```
syntax on
set nu
set ruler
set autoindent
set smartindent
set expandtab
set tabstop=4
set shiftwidth=4
```

### 1.2 Starter Code

```
#include <bits/stdc++.h>
#define LL long long
#define F first
#define S second

using namespace std;

int main () {
   cin.tie(0);
   ios_base::sync_with_stdio(0);

   return 0;
}
```

## 1.3 C++ Grammar, STL

```
struct Cmp { bool operator() (T &a, T &b) { return true; } };
set<T,Cmp> s;
bool cmp (T &a, T &b) { return true; }
set<T,decltype(cmp)> s(cmp);
auto cmp = [](T &a, T &b) -> bool { return true; }
set<T,decltype(cmp)> s(cmp);

map<int,int> m;
m.find(val) == m.end()
for (auto p : m) { key = p.F; value = p.S; }

priority_queue<T, vector<T>, Cmp> pq;
```

## 2 Data Structures

### 2.1 Binary Indexed Tree 2D

```
// Support 2 types of queries:
// - Add v to cell (x, y)
// - Get the sum of rectangle with top-left corner (1, 1)
// and lower-right corner (x, y)
void update(int x, int y, int v) {
   while (x \le n) {
       int z = y;
       while (z \le n) \{
           bit[x][z] += v;
           z += (z \& (-z));
       x += (x & (-x));
int get(int x, int y) {
   if (x == 0 || y == 0) return 0;
   int sum = 0;
   while (x) {
       int z = y;
       while (z) {
           sum += bit[x][z];
           z = (z \& (-z));
       x = (x & (-x));
```

```
}
return sum;
}
```

## 2.2 Segment Tree 2D

```
// Supported:
// - Add a value v to cell (x, y)
// - Get the sum in rectangle with top left corner
// (x1, y1) and bottom right corner (x2, y2)
void build_y(int k_x, int k_y, int 1, int r) {
   if (1 == r) {
       t[k_x][k_y] = 0;
       return;
   }
   int mid = (1 + r) >> 1;
   build_y(k_x, k_y * 2, 1, mid);
   build_y(k_x, k_y * 2 + 1, mid + 1, r);
   t[k_x][k_y] = 0;
}
void build_x(int k, int l, int r) {
   build_y(k, 1, 1, n);
   if (1 == r) return;
   int mid = (1 + r) >> 1;
   build_x(k * 2, 1, mid);
   build_x(k * 2 + 1, mid + 1, r);
}
void update_y(int k_x, int l_x, int r_x, int k_y, int l_y, int r_y,
    int v, int v) {
   if (y < l_y || r_y < y) return;
   if (1_y == r_y) {
       if (1_x == r_x)
          t[k_x][k_y] += v;
       else
           t[k_x][k_y] = t[k_x * 2][k_y] + t[k_x * 2 + 1][k_y];
       return;
   int mid = (l_y + r_y) >> 1;
   update_v(k_x, 1_x, r_x, k_y * 2, 1_y, mid, y, v);
   update_y(k_x, l_x, r_x, k_y * 2 + 1, mid + 1, r_y, y, v);
   t[k_x][k_y] = t[k_x][k_y * 2] + t[k_x][k_y * 2 + 1];
```

```
}
void update_x(int k, int l, int r, int x, int y, int v) {
    if (x < 1 \mid | r < x) return;
    if (1 == r) {
       update_y(k, l, r, 1, 1, n, y, v);
       return;
    }
    int mid = (1 + r) >> 1;
    update_x(k * 2, 1, mid, x, y, v);
    update_x(k * 2 + 1, mid + 1, r, x, y, v);
    update_y(k, l, r, 1, 1, n, y, v);
int get_y(int k_x, int k_y, int 1, int r, int y1, int y2) {
    if (y2 < 1 || r < y1) return 0;
    if (y1 \le 1 \&\& r \le y2) return t[k_x][k_y];
    int mid = (1 + r) >> 1;
    return get_y(k_x, k_y * 2, 1, mid, y1, y2) +
          get_y(k_x, k_y * 2 + 1, mid + 1, r, y1, y2);
}
int get_x(int k, int l, int r, int x1, int x2, int y1, int y2) {
   if (r < x1 || x2 < 1) return 0;
    if (x1 <= 1 && r <= x2)
       return get_y(k, 1, 1, n, y1, y2);
    int mid = (1 + r) >> 1;
    return get_x(k * 2, 1, mid, x1, x2, y1, y2) +
          get_x(k * 2 + 1, mid + 1, r, x1, x2, y1, y2);
}
```

## 2.3 Persistent Segment Tree

```
struct Node {
   Node() = default;

Node(int 1, int r, int v)
      : left(l), right(r), val(v) {}

int left, right, val;
};

int build(int k, int 1, int r) {
   tree[k].val = 0;
```

```
if (1 == r) return k;
   tree[k].left = ++num_node;
   tree[k].right = ++num_node;
   int mid = (1 + r) >> 1;
   build(tree[k].left, 1, mid);
   build(tree[k].right, mid + 1, r);
   return k;
}
int update(int k, int l, int r, int i, int v) {
   int K = ++num_node;
   if (1 == r) {
       tree[K].val = tree[k].val + v;
       return K;
   }
   tree[K].left = tree[k].left;
   tree[K].right = tree[k].right;
   int mid = (1 + r) >> 1;
   if (i <= mid)</pre>
       tree[K].left = update(tree[K].left, 1, mid, i, v);
   else
       tree[K].right = update(tree[K].right, mid + 1, r, i, v);
   tree[K].val = tree[tree[K].left].val + tree[tree[K].right].val;
   return K;
}
```

## 2.4 Treap

```
// Treap
// Tested on POJ 2828

struct Node {
    int v, k, size, cnt;
    Node *1, *r;
    Node (int v) : v(v), k(rand()), size(1), cnt(1), l(NULL), r(NULL)
        {}
    void update () {
        size = (1 ? 1->size : 0) + (r ? r->size : 0) + cnt;
    }
};
Node *root;

void zig (Node* &u) {
    Node *v = u->1;
```

```
if (!v) return;
    u->1 = v->r; v->r = u;
   v->size = u->size; u->update();
    u = v:
}
void zag (Node* &u) {
   Node *v = u->r;
    if (!v) return;
   u->r = v->1; v->1 = u;
    v->size = u->size; u->update();
    u = v:
}
void insert (Node* &u, int v) {
    if (!u) { u = new Node(v); return; }
    if (v < u \rightarrow v) {
       insert(u->1, v);
       if (u->l->k < u->k) zig(u);
   } else if (v > u -> v) {
       insert(u->r, v);
       if (u->r->k < u->k) zag(u);
   } else {
       u->cnt++;
    }
    u->update();
```

## 2.5 Splay Tree

```
// Supports reversing a segment.

struct SplayTree {
    struct Node {
        Node *left, *right, *parent;
        int value, size;
        bool reversed;
    };

SplayTree() {
        nilt = new Node();
        nilt->left = nilt->right = nilt->parent = nilt;
        nilt->value = nilt->size = 0;
        nilt->reversed = false;
```

```
}
void set_left(Node* x, Node* y) {
   x \rightarrow left = y;
   y->parent = x;
}
void set_right(Node* x, Node* y) {
   x->right = y;
   y->parent = x;
}
void set_child(Node* x, Node* y, bool is_right) {
    if (is_right) set_right(x, y);
    else set_left(x, y);
}
void build_tree(vector < int >& arr) {
   root = nilt;
   for (int i = 0; i < arr.size(); ++i) {
       Node* x = new Node();
       x->size = arr[i];
       x->value = arr[i];
       x->reversed = false;
       set_left(x, root);
       x->parent = x->right = nilt;
       root = x;
   }
}
void propagate(Node* x) {
   if (x == nilt) return;
    if (x->reversed) {
       swap(x->left, x->right);
       x->left->reversed = !x->left->reversed;
       x->right->reversed = !x->right->reversed;
       x->reversed = false;
   }
}
Node* locate(Node* x, int pos) {
   do {
       propagate(x);
       int num = x->left->size + 1;
       if (num == pos) return x;
       if (num > pos) x = x -> left;
```

```
else
           pos -= num, x = x->right;
   } while (true);
   return x:
}
void update(Node* x) {
   x->size = x->left->size + x->right->size + 1;
}
void uptree(Node* x) {
   Node* y = x->parent;
   Node* z = y->parent;
   if (x == y->right) {
       Node* b = x->left;
       set_right(y, b);
       set_left(x, y);
   }
   else {
       Node* b = x->right;
       set_left(y, b);
       set_right(x, y);
   update(v);
   update(x);
   set_child(z, x, z->right == y);
}
void splay(Node* x) {
   do {
       Node* y = x-parent;
       if (y == nilt) return;
       Node* z = y->parent;
       if (z != nilt) {
           if ((x == y -> left) == (y == z -> left))
              uptree(y);
           else
              uptree(x);
       }
       uptree(x);
   } while (true);
void split(Node* t, int pos, Node*& t1, Node*& t2) {
   if (pos == 0) {
       t1 = nilt;
```

```
t2 = t:
           return;
       }
       if (pos \geq t-\geqsize) {
           t1 = t;
           t2 = nilt;
           return;
       }
       Node* x = locate(t, pos);
       splay(x);
       t1 = x;
       t2 = x->right;
       t1->right = nilt;
       t2->parent = nilt;
       update(t1);
   }
   Node* join(Node* t1, Node* t2) {
       if (t1 == nilt) return t2;
       t1 = locate(t1, t1->size);
       splay(t1);
       set_right(t1, t2);
       update(t1);
       return t1;
   }
   Node *root, *nilt;
};
```

## 2.6 Mo's Algorithm

## 3 Graph Theory

#### 3.1 Ford Fulkerson

```
bool find_path() {
   int l = 1, r = 1; ++flag;
   q[1] = source; check[source] = flag;
   while (1 <= r) {
       int u = q[1++];
       for (auto v : adj[u])
           if (check[v] != flag && c[u][v] > f[u][v]) {
              pre[v] = u;
              check[v] = flag;
              if (v == target) return true;
              q[++r] = v;
           }
   }
   return false;
}
void augment() {
   int v = target, delta = oo;
   while (v != source) {
       int u = pre[v];
       delta = min(c[u][v] - f[u][v], delta);
       v = u;
   v = target; flow += delta;
   while (v != source) {
       int u = pre[v];
       f[u][v] += delta;
       f[v][u] -= delta:
       v = u;
   }
}
```

## 3.2 Dinic

```
const int MAXN = 1024;

struct Edge {
  int u, v, c;
  Edge *next, *rev;
  void set(int u, int v, int c, Edge *next, Edge *rev) {
    this->u = u;
    this->v = v;
    this->c = c;
}
```

```
this->next = next:
   this->rev = rev;
};
struct Node {
 Edge *head;
 int level;
 Node() : head(NULL), level(-1) {}
};
struct Graph {
 int n, m;
 Node *nodes;
 Edge *edges;
 Graph() {
   cin >> n >> m;
   nodes = new Node[n];
   edges = new Edge[2*m];
   for (int i = 0; i < m; i ++) {
     int u, v, c;
     cin >> u >> v >> c;
     edges[2*i].set(u, v, c, nodes[u].head, &edges[2*i+1]);
     nodes[u].head = &edges[2*i];
     edges[2*i+1].set(v, u, 0, nodes[v].head, &edges[2*i]);
     nodes[v].head = &edges[2*i+1];
   }
 }
 bool make_level() {
   for (int i = 0; i < n; i ++) {
     nodes[i].level = -1;
   }
   queue<Node*> queue;
   queue.push(&nodes[0]);
   nodes[0].level = 0;
   while (!queue.empty()) {
     Node* node = queue.front();
     queue.pop();
     for (Edge *edge = node->head; edge; edge = edge->next) {
       if (nodes[edge->v].level == -1 && edge->c) {
         nodes[edge->v].level = node->level + 1;
         queue.push(&nodes[edge->v]);
     }
```

```
}
   return nodes[n-1].level != -1;
 }
  int find(int u, int key) {
   if (u == n-1) return key;
   for (Edge *edge = nodes[u].head; edge; edge = edge->next) {
     if (nodes[edge->v].level == nodes[u].level + 1 && edge->c) {
       int flow = find(edge->v, min(key, edge->c));
       if (flow) {
         edge->c -= flow;
         edge->rev->c += flow;
         return flow;
       }
     }
   return 0;
  int dinic() {
   int ans = 0;
   int flow;
   while (make_level()) {
     while ((flow = find(0, INT_MAX))) {
       ans += flow;
     }
   }
   return ans;
 }
};
```

## 3.3 Tarjan

```
// avail[] initialized to be all 0
void tarjan(int u) {
   num[u] = low[u] = ++num_node;
   st.push(u);
   for (int i = 0; i < adj[u].size(); ++i) {
      int v = adj[u][i];
      if (!avail[v]) {
        if (num[v] == 0) {
            tarjan(v);
            low[u] = min(low[u], low[v]);
      }
}</pre>
```

```
else low[u] = min(low[u], num[v]);
}
if (low[u] == num[u]) {
   int v = -1;
   ++num_comp;
   while (v != u) {
      v = st.top(); st.pop();
      comp[v] = num_comp;
      avail[v] = 1;
}
}
```

## 3.4 Topo Sort

```
void topo_sort() {
   for (int i = 1; i <= num_comp; ++i)
      if (deg[i] == 0) q.push(i);
   int num = 0;
   while (!q.empty()) {
      int u = q.front(); q.pop();
      for (int i = 0; i < new_adj[u].size(); ++i) {
        int v = new_adj[u][i];
        --deg[v];
      if (deg[v] == 0) q.push(v);
      }
      position[u] = ++num;
   }
}</pre>
```

### 3.5 2-SAT

```
++deg[comp[v]];
}

topo_sort();
for (int i = 0; i < list_node.size(); ++i) {
   int u = list_node[i];
   // position[u]: position of u after topo sorted
   if (position[comp[u]] > position[comp[neg[u]]])
        check[u] = 1; // Pick u (otherwise pick !u)
}
return true;
```

## 3.6 Lowest Common Ancestor - $O(n \log n)$

```
// Note: Log = ceil(log2(n))
// d[u] = depth of node u + 1 (ie: d[root] = 1)
void buildLCA() {
   for (int i = 1; i <= n; ++i) p[i][0] = par[i];
   for (int j = 1; j \le Log; ++j)
       for (int i = 1; i \le n; ++i)
           p[i][j] = p[p[i][j-1]][j-1];
}
int LCA(int u, int v) {
   if (d[u] < d[v]) swap(u, v);
   for (int j = Log; j \ge 0; --j)
       if (d[p[u][j]] >= d[v]) u = p[u][j];
   if (u == v) return u;
   for (int j = Log; j \ge 0; --j)
       if (p[u][j] != p[v][j]) {
           u = p[u][j];
           v = p[v][j];
   return p[u][0];
```

### 3.7 Eulerian Circuit

```
// adj[] is unordered_map
void euler(int start) {
   stack < int > st; st.push(start);
```

#### 3.8 Dominator Tree

```
// For multitest, initialize label[], l[], r[] to 0 and
// clear adj[], radj[], bucket[], child[].
struct DominatorTree {
   DominatorTree() = default;
   bool reachable(int u) { return label[u] > 0; }
   void add_edge(int u, int v) { adj[u].push_back(v); }
   void dfs(int u) {
       label[u] = ++num; orig[num] = u;
       dsu[num] = mdsu[num] = sdom[num] = num;
       for (int v : adj[u]) {
          if (!label[v]) {
              dfs(v);
              parent[label[v]] = label[u];
          radj[label[v]].push_back(label[u]);
   }
   int find_dsu(int u, int flag = 0) {
       if (u == dsu[u]) return flag ? -1 : u;
       int p = find_dsu(dsu[u], 1);
       if (p < 0) return u;
       if (sdom[mdsu[dsu[u]]] < sdom[mdsu[u]])</pre>
```

```
mdsu[u] = mdsu[dsu[u]];
       dsu[u] = p;
       return flag ? p : mdsu[u];
   void dfs_dominator(int u) {
       l[u] = ++num:
       for (int v : child[u]) dfs_dominator(v);
       r[u] = num:
   bool dominates(int u. int v) {
       return l[u] <= l[v] && l[v] <= r[u];
   }
   void build() {
       num = 0; dfs(root);
       for (int u = num; u >= 1; --u) {
           for (int v : radj[u])
              sdom[u] = min(sdom[u], sdom[find_dsu(v)]);
           if (u != 1) bucket[sdom[u]].push_back(u);
           for (int v : bucket[u]) {
              int w = find_dsu(v);
              if (sdom[w] == sdom[v]) idom[v] = sdom[v];
              else idom[v] = w;
           if (u != 1) dsu[u] = parent[u];
       for (int i = 2; i \le num; ++i) {
           if (idom[i] != sdom[i]) idom[i] = idom[idom[i]];
           child[orig[idom[i]]].push_back(orig[i]);
       }
       num = 0; dfs_dominator(root);
   }
   vector < int > adj[maxN], radj[maxN], bucket[maxN], child[maxN];
   int sdom[maxN], idom[maxN], dsu[maxN], mdsu[maxN], n, root, num;
   int parent[maxN], label[maxN], orig[maxN], l[maxN], r[maxN];
};
```

## 3.9 Centroid Decomposition

```
void build(int u, int p) {
   sze[u] = 1;
```

```
for (int v : adj[u])
       if (!elim[v] && v != p) build(v, u), sze[u] += sze[v];
}
int get_centroid(int u, int p, int num) {
   for (int v : adj[u])
       if (!elim[v] && v != p && sze[v] > num / 2)
           return get_centroid(v, u, num);
    return u;
void centroid_decomposition(int u) {
   build(u, -1);
   int root = get_centroid(u, -1, sze[u]);
   // Do stuffs here
   elim[root] = true;
   for (int v : adj[root])
       if (!elim[v]) centroid_decomposition(v, c + 1);
}
```

## 3.10 Heavy Light Decomposition

```
void build(int u) {
    size tree[u] = 1:
   for (int i = 0; i < adj[u].size(); ++i) {
       int v = adj[u][i];
       if (parent[u] == v) continue;
       parent[v] = u;
       build(v);
       size_tree[u] += size_tree[v];
}
void hld(int u) {
   if (chain_head[num_chain] == 0)
       chain_head[num_chain] = u;
    chain_idx[u] = num_chain;
    arr_idx[u] = ++num_arr;
   node_arr[num_arr] = u;
   int heavy_child = -1;
   for (int i = 0; i < adj[u].size(); ++i) {
       int v = adj[u][i];
       if (parent[u] == v) continue;
```

```
if (heavy_child == -1 || size_tree[v] > size_tree[heavy_child])
           heavy_child = v;
   }
   if (heavy_child != -1)
       hld(heavy_child);
   for (int i = 0; i < adj[u].size(); ++i) {</pre>
       int v = adj[u][i];
       if (v == heavy_child || parent[u] == v) continue;
       ++num_chain;
       hld(v);
}
// u is an ancestor of v
int query_hld(int u, int v) {
   int uchain = chain_idx[u], vchain = chain_idx[v], ans = -1;
   while (true) {
       if (uchain == vchain) {
           get(..., arr_idx[u], arr_idx[v]);
           break;
       get(..., arr_idx[chain_head[vchain]], arr_idx[v]);
       v = parent[chain_head[vchain]];
       vchain = chain_idx[v];
   }
   return ans;
```

## 4 Dynamic Programming

### 4.1 Convex Hull Trick

```
}
// Assuming lines' slopes m are strictly increasing.
void add(htype m, htype b) {
    while (lst.size() >= 2 && is_bad(lst[lst.size() - 2], lst.back(),
       {m, b}))
       lst.pop_back();
   lst.push_back({m, b});
}
htype get_value(line d, htype x) {
   return d.first * x + d.second:
}
// Assuming queries' x are strictly increasing.
int pointer = 0;
htype get(htype x) {
   if (pointer > lst.size()) pointer = lst.size() - 1;
   while (pointer < lst.size() - 1 && get_value(lst[pointer], x) <</pre>
        get_value(lst[pointer + 1], x))
       ++pointer;
   return get_value(lst[pointer], x);
}
```

## 4.2 Dynamic Convex Hull Trick

```
// Slow but correct. Takes O(log n) per add and query.

typedef long long htype;

// Representing a line. To query value x,

// set m = x, is_query = true.

struct Line {
    bool operator < (const Line& rhs) const {
        // Compare lines
        if (!rhs.is_query) return m < rhs.m;

        // Compare queries
        const Line* s = nxt();
        if (s == NULL) return false;
        htype x = rhs.m;
        return s->m * x + s->b > m * x + b;
    }

    htype m, b;
```

```
bool is_query;
    mutable function < const Line*() > nxt;
};
class ConvexHullTrick : public set < Line > {
  public:
    void add(htype m, htype b) {
       auto p = insert({m, b, false});
       if (!p.second) return;
       iterator y = p.first;
       y->nxt = [=] { return (next(y) == end()) ? NULL : &(*next(y));
           };
       if (is_bad(v)) {
           erase(y);
           return;
       while (next(y) != end() && is_bad(next(y))) erase(next(y));
       while (y != begin() && is_bad(prev(y))) erase(prev(y));
   }
    htype get(htype x) {
       iterator y = lower_bound({x, 0, true});
       return y->m * x + y->b;
   }
  private:
   bool is_bad(iterator y) {
       iterator z = next(y);
       if (y == begin())
           return ((z == end()) ? false : y->m == z->m && y->b <=
               z->b);
       iterator x = prev(y);
       if (z == end())
           return (y->m == x->m \&\& y->b <= x->b);
       return (x-b - y-b) * (z-m - y-m) >= (y-b - z-b) * (y-m - y-m)
           x->m):
   }
};
```

## 5 String

#### 5.1 Z-Function

```
// z[] is 1-based, z[1] = 0
```

### 5.2 Suffix Array

```
bool suffix_cmp(int i, int j) {
   if (pos[i] != pos[j]) return pos[i] < pos[j];</pre>
   i += gap;
   j += gap;
   return (i < N && j < N) ? pos[i] < pos[j] : i > j;
}
void build sa() {
   N = s.size();
   for (int i = 0; i < N; ++i) sa[i] = i, pos[i] = s[i];
   for (gap = 1;; gap *= 2) {
       sort(sa, sa + N, suffix_cmp);
       for (int i = 0; i < N - 1; ++i) tmp[i + 1] = tmp[i] +
           suffix_cmp(sa[i], sa[i + 1]);
       for (int i = 0; i < N; ++i) pos[sa[i]] = tmp[i];</pre>
       if (tmp[N-1] == N-1) break;
}
// height[i] = length of common prefix of suffix(sa[i]) and
    suffix(sa[i+1])
void build_height () {
   height.assign(n-1, -1);
   for (int i = 0, k = 0; i < n; i++) {
       if (rk[i] == n-1) continue;
       if (k) k--;
       for (int j = sa[rk[i]+1]; i+k< n && j+k< n && s[i+k] == s[j+k];
           k++);
```

```
height[rk[i]] = k;
}
```

#### 5.3 Aho-Corasick Automata

```
struct Node {
 Node* next[26]:
 Node* fail;
 int cnt;
 Node (Node* root) {
   memset(next, NULL, sizeof(next));
   fail = root;
   cnt = 0;
 }
};
Node* root;
void insert (string s) {
 Node* curr = root;
 for (int i = 0; i < s.length(); i++) {
   int j = s[i] - 'a';
   if (curr->next[j] == NULL) {
     curr->next[j] = new Node(root);
   curr = curr->next[j];
 curr->cnt++;
void make_fail () {
 queue<Node*> q;
 for (int i = 0; i < 26; i++) {
   if (root->next[i]) {
     q.push(root->next[i]);
 while (!q.empty()) {
   Node* node = q.front(); q.pop();
   for (int i = 0; i < 26; i++) {
     if (node->next[i]) {
       q.push(node->next[i]);
       Node* f = node->fail;
       while (f != root && !f->next[i]) {
```

```
f = f->fail:
       }
       if (f->next[i]) {
         f = f \rightarrow next[i];
       node->next[i]->fail = f;
   }
 }
int work (string s) {
 set<Node*> seen;
 int cnt = 0;
 Node* curr = root;
 for (int i = 0; i < s.length(); i++) {</pre>
   int j = s[i] - 'a';
   while (curr != root && !curr->next[j]) {
     curr = curr->fail;
   }
   if (curr->next[j]) {
     curr = curr->next[j];
     Node* p = curr;
     while (p != root) {
       if (seen.find(p) != seen.end()) break;
       seen.insert(p);
       cnt += p->cnt;
       p = p->fail;
   }
 }
 return cnt;
```

### 5.4 Palindromic Tree

```
struct Node {
   Node* next[26]; // to palindrome by extending me with a letter
   Node* sufflink; // my LSP
   int len; // length of this palindrome substring
   int num; // number of palindrome substrs ending here
};
Node nodes[NMAX];
int n = 0; // number of nodes in tree
```

```
vector<int> s;
LL ans = 0;
void build tree () {
    nodes[0].len = -1; nodes[0].sufflink = &nodes[0]; // root 0
    nodes[1].len = 0; nodes[1].sufflink = &nodes[0]; // root 1
    n = 2;
    Node* suff = &nodes[1]; // node for LSP of processed prefix
    for (int i = 0; i < s.size(); i++) {
       // find LSP xAx
       Node* ptr = suff;
       while (1) {
           int j = i - 1 - ptr \rightarrow len;
           if (j \ge 0 \&\& s[j] == s[i]) break;
           ptr = ptr->sufflink;
       }
       if (ptr->next[s[i]]) { // palindrome substr already exists
           suff = ptr->next[s[i]];
       } else { // add a new node
           suff = &nodes[n++];
           suff->len = ptr->len + 2;
           ptr->next[s[i]] = suff;
           if (suff->len == 1) { // current LSP is trivial
               suff->sufflink = &nodes[1];
               suff \rightarrow num = 1;
           } else {
               // find xAx's LSP xBx
               while (1) {
                   ptr = ptr->sufflink;
                   int j = i - 1 - ptr \rightarrow len;
                   if (i \ge 0 \&\& s[i] == s[i]) break;
               suff->sufflink = ptr->next[s[i]];
               suff->num = suff->sufflink->num + 1;
           }
       }
       ans += suff->num:
    }
```

## 6 Game Theory

#### 6.1 Nim Product

```
// Note: (i | j) might overflow
int nim_multiply(int x, int y) {
   int p = 0;
   for (int i = 0; i < maxLog + 1; ++i)
       if (x & (1 << i))
           for (int j = 0; j < maxLog + 1; ++j)
               if (y & (1 << j))
                  p ^= mul[i][j];
}
void init() {
   for (int i = 0; i < maxLog + 1; ++i)
       for (int j = 0; j \le i; ++j) {
           if ((i \& j) == 0) \text{ mul}[i][j] = 1 << (i | j);
               mul[i][j] = 1;
               for (int t = 0; t < \max Log + 1; ++t) {
                  int k = (1 << t);
                   if (i & j & k) mul[i][j] = nim_multiply(mul[i][j],
                       ((1 << k) * 3) >> 1);
                  else
                      if ((i | j) & k) mul[i][j] =
                           nim_multiply(mul[i][j], (1 << k));</pre>
               }
           mul[j][i] = mul[i][j];
}
```

## 7 Math

## 7.1 Number Theory

```
long long gcd (long long a, long long b) { return b == 0 ? a : gcd(b,
    a%b); }
long long mul_mod (long long x, long long y, long long MOD) {
```

```
long long q = (long long)((long double)x * y / MOD);
long long r = x * y - q * MOD;
while (r < 0) r += MOD;
while (r >= MOD) r -= MOD;
return r;
}
long long pow_mod (long long b, long long e, long long MOD) {
  long long ans = 1;
  while (e) {
    if (e & 1) ans = mul_mod(ans, b, MOD);
        b = mul_mod(b, b, MOD);
        e >>= 1;
    }
return ans;
}
```

#### 7.1.1 Extended Euclid

```
// Extended Euclid
// Solve xa + yb = gcd(a, b)
pair<long long,pair<long long,long long>> extended_euclid (long long
        a, long long b) {
    if (b == 0) return {a, {1, 0}};
    auto ee = extended_euclid(b, a % b);
    long long g = ee.first;
    long long y = ee.second.first;
    long long x = ee.second.second;
    y -= a / b * x;
    return {g, {x, y}};
}
```

#### 7.1.2 Mod Linear Equation

```
x %= n / g; x += n / g; x %= n / g;
return x;
}
```

#### 7.1.3 Chinese Remainder Theorem

```
// Chinese Remainder Theorem
// Solve x = bi (mod mi)
long long chinese_remainder_theorem (vector<long long> b, vector<long
    long> m) {
    int n = b.size();
    long long M = 1, ans = 0;
    for (int i = 0; i < n; i++) M *= m[i];
    for (int i = 0; i < n; i++) {
        long long Mi = M / m[i];
        auto ee = extended_euclid(Mi, m[i]);
        long long xi = ee.second.first;
        ans += Mi * xi * b[i];
    }
    ans %= M; ans += M; ans %= M;
    return ans;
}</pre>
```

### 7.1.4 Miller-Rabin prime test

```
// Miller-Rabin prime test O(log(n)^3)
// Tested on UVA 11476
bool miller_rabin (long long n, long long a) {
 if (n == 2 \mid \mid n == a) return true;
 if ((n \& 1) == 0) return false:
 int s = 0; long long d = n - 1; while (!(d & 1)) { d >>= 1; s++; }
 long long t = pow_mod(a, d, n);
 if (t == 1 || t == n-1) return true;
 for (; s; s--) {
   t = mul_mod(t, t, n);
   if (t == n-1) return true;
 }
 return false;
bool is_prime (long long n) {
 if (n < 2) return false;
 vector<int> va = {2,3,5,7,11,13,17,19,23,29,31,37};
 for (int a : va) {
```

```
if (!miller_rabin(n, a)) return false;
}
return true;
```

#### 7.1.5 Pollard rho prime factorization

```
// Pollard rho prime factorization O(n^{0.25})
// Tested on UVA 11476
long long pollard_rho (long long n) {
 // find a non-trivial prime factor of n
 // n must not be a prime (will loop forever!)
 while (1) {
   long long c = rand() \% (n-1) + 1;
   long long x, y; x = y = rand() \% (n-1) + 1;
   long long head = 1, tail = 2;
   while (1) {
     x = (mul_mod(x, x, n) + c) \% n;
     if (x == v) break:
     auto d = gcd(abs(x-y), n);
     if (d > 1 && d < n) return d;
     if ((++head) == tail) { y = x; tail <<= 1; }</pre>
 }
map<long long,int> factorize (long long n) {
 if (n == 1) return {};
 if (is_prime(n)) return {{n, 1}};
 map<long long,int> fac;
 auto p = pollard_rho(n);
 auto fac0 = factorize(p);
 auto fac1 = factorize(n/p);
 for (auto be : fac0) fac[be.first] += be.second;
 for (auto be : fac1) fac[be.first] += be.second:
 return fac;
```

#### 7.1.6 Primitive root

```
// Primitive root
// p is prime
long long primitive_root (long long p) {
  auto fac = factorize(p - 1);
```

```
for (long long g = 1; ; g++) {
   bool ok = true;
   for (auto be : fac) {
      long long b = be.first;
      if (pow_mod(g, (p - 1) / b, p) == 1) { ok = false; break; }
    }
   if (ok) return g;
}
return -1; // should never reach here
}
```

#### 7.1.7 Discrete log

```
// Discrete log O(p^0.5)
// Solve a^x = b (mod p) (p is prime)
long long discrete_log (long long a, long long b, long long p) {
 long long rp = (long long)sqrt(p);
 map<long long,long long> rec;
 long long tmp = 1:
 for (long long i = 0; i < rp; i++) {
   rec[tmp] = i;
   tmp = tmp * a % p;
 int cur = 1:
 for (long long q = 0; q*rp < p; q++) {
   long long r = mod_linear_equation(cur, b, p);
   if (rec.find(r) != rec.end()) return q * rp + rec[r];
   cur = cur * tmp % p;
 return -1; // no solution
}
```

#### 7.1.8 Exp remainder

```
// Exp remainder O(p^0.5)
// Solve x^a = b (mod p) (p is prime)
long long exp_remainder (long long a, long long b, long long p) {
  long long g = primitive_root(p);
  long long s = discrete_log(g, b, p);
  if (b == 0) return 0;
  if (s == -1) return -1;
  auto fac = extended_euclid(a, p-1);
  long long d = fac.first;
```

```
long long x = fac.second.first;
long long y = fac.second.second;
if (s % d) return -1;
x = x * s/d;
x %= p-1; x += p-1; x %= p-1;
for (long long i = 0; i < d; i++) x = (x + (p-1)/d) % (p-1);
return pow_mod(g, x, p);
}</pre>
```

#### 7.1.9 Euler function

```
// Euler function O(n^0.5)
long long phi (long long n, long long key = 2) {
 if (n == 1) return 1;
 while (n % key && key * key \leq n) key++;
 if (key * key > n) return n-1;
 if (n / key % key) return phi(n/key, key+1) * (key-1);
 return phi(n/key, key) * key;
// Euler function preprocess O(nlogn)
void phi_gen (int n) {
 vector<int> mindiv(n+1, 0), phi(n+1, 0);
 for (int i = 1; i <= n; i++) mindiv[i] = i;
 for (int i = 2: i*i <= n: i++) {
   if (mindiv[i] != i) continue;
   for (int j = i*i; j <= n; j += i) mindiv[j] = i;
 phi[1] = 1;
 for (int i = 2; i <= n; i++) {
   phi[i] = phi[i / mindiv[i]];
   if ((i / mindiv[i]) % mindiv[i] == 0) phi[i] *= mindiv[i];
   else phi[i] *= mindiv[i] - 1;
 }
```

#### 7.1.10 Mobiüs function

```
// Mobius function O(n^0.5)
long long mu (long long n) {
  auto fac = factorize(n);
  for (auto be : fac) {
    if (be.second > 1) return 0;
  }
```

```
return (fac.size() % 2 == 0) ? 1 : -1;
}
// Mobius function preprocess O(nlogn)
void mu_gen (int n) {
  vector<int> mu(n+1, 0);
  for (int i = 1; i <= n; i++) {
    int target = i == 1;
    int delta = target - mu[i];
    mu[i] = delta;
    for (int j = i+i; j <= n; j += i) mu[j] += delta;
  }
}</pre>
```

### 7.2 Matrix

#### 7.2.1 Matrix inverse

#### 7.2.2 rref

#### 7.2.3 Gaussian Elimination

```
// Note: ax = b
bool gaussian_elimination() {
   vector < int > row;
   for (int i = 0; i < N; ++i) row.push_back(i);
   for (int t = 0; t < N; ++t) {
       int R = -1;
       for (int i = t; i < N; ++i) {
           int r = row[i];
          if (a[r][t] > eps) {
              R = i;
              break:
          }
       }
       if (R == -1) return false;
       swap(row[R], row[t]);
       R = row[t];
       for (int i = t + 1; i < N; ++i) {
           int r = row[i];
           double p = a[r][t] / a[R][t];
          for (int c = 0; c < N; ++c)
              a[r][c] -= p * a[R][c];
          b[r] -= p * b[R];
       }
   }
```

```
for (int i = N - 1; i >= 0; --i) {
    int r = row[i];
    for (int c = N - 1; c > i; --c)
        b[r] -= a[r][c] * res[c];
    res[r] = b[r] / a[r][i];
}
return true;
}
```

#### 7.3 Discrete Fourier Transform

#### 7.3.1 Base Class

```
// To multiply a, b and put result in c:
// PolyMul::polynomial_multiply(a, b, c);
template < class Transform >
struct DFT {
   #define TAdd Transform::add
   #define TSub Transform::subtract
   #define TMul Transform::multiply
   typedef vector < int64_t > ivector;
   typedef typename Transform::ctype DType;
   typedef vector < DType > dvector;
   typedef vector < vector < dvector > > mdvector:
   static void init() {
       w.resize(NBIT);
       for (int iter = 0, len = 1; iter < NBIT; ++iter, len *= 2) {
           w[iter].resize(2);
           for (int invert = 0; invert < 2; ++invert) {</pre>
              w[iter][invert].assign(1 << iter, 0);
              DType wlen = Transform::generate_root(2 * len, invert);
              w[iter][invert][0] = 1;
              for (int j = 1; j < len; ++j)
                  w[iter][invert][j] = TMul(w[iter][invert][j - 1],
                      wlen);
          }
       }
   static void fft(dvector& a, bool invert = false) {
       int n = a.size();
       for (int i = 1, j = 0; i < n; ++i) {
           int bit = n >> 1;
           for (; j & bit; bit >>= 1) j ^= bit;
           j ^= bit;
```

```
if (j > i) swap(a[i], a[j]);
       }
       for (int iter = 0, len = 1; len < n; ++iter, len *= 2) {
           DType wlen = Transform::generate_root(2 * len, invert);
           for (int i = 0; i < n; i += 2 * len) {
              for (int j = 0; j < len; ++j) {
                  auto x = a[i + j];
                  auto y = TMul(w[iter][invert][j], a[i + j + len]);
                  a[i + j] = TAdd(x, y);
                  a[i + j + len] = TSub(x, y);
              }
          }
       if (invert) Transform::invert(a);
   }
   static void polynomial_multiply(
       const ivector& a, const ivector& b, ivector& out) {
       uint32_t new_size = a.size() + b.size() - 1;
       for (NBIT = 0, N = 1; N < new_size; N *= 2, ++NBIT) {}
       dvector fa(a.begin(), a.end()), fb(b.begin(), b.end());
       fa.resize(N); fft(fa);
       fb.resize(N); fft(fb);
       for (int i = 0; i < fa.size(); ++i) fa[i] = TMul(fa[i], fb[i]);
       fft(fa, true);
       Transform::prepare_output(fa, out, new_size);
   static int32_t NBIT, N;
   static mdvector w:
};
// Remember to call PolyMul::init() in main().
using PolyMul = DFT < FFT >;
template<> int32_t PolyMul::NBIT = /* max of log(n) */;
template<> int32_t PolyMul::N = 1 << PolyMul::NBIT;</pre>
template<> PolyMul::mdvector PolyMul::w = PolyMul::mdvector();
```

#### 7.3.2 Fast Fourier Transform

```
struct FFT {
   typedef vector < int64_t > ivector;
   typedef complex < double > ctype;
   typedef vector < ctype > cvector;
   static ctype add(ctype x, ctype y) { return x + y; }
   static ctype subtract(ctype x, ctype y) { return x - y; }
   static ctype multiply(ctype x, ctype y) { return x * y; }
```

```
static ctype generate_root(int len, bool invert) {
       double alpha = 2.0 * PI / len * (invert ? -1 : 1);
       return ctype(cos(alpha), sin(alpha));
   static void prepare_output(
       const cvector& vin, ivector& vout, uint32_t out_size) {
       vout.resize(out_size);
       for (int i = 0; i < out_size; ++i)</pre>
           vout[i] = llround(vin[i].real());
       while (vout.size() > 1 && vout.back() == 0)
           vout.pop_back();
   }
   static void invert(cvector& a) {
       for (auto\& x : a) x /= a.size();
   static double PI;
};
double FFT::PI = acos(-1.0);
```

#### 7.3.3 Number Theoretic Transform

```
struct NTT {
   typedef vector < int64_t > ivector;
   typedef int64_t ctype;
   typedef vector < ctype > cvector;
   static ctype add(ctype x, ctype y) {
       return 111 * x + y < mod ? x + y : x + y - mod;
   static ctype subtract(ctype x, ctype y) {
       return x < y? 111 * x - y + mod : x - y;
   static ctype multiply(ctype x, ctype y) {
       return (111 * x * y) % mod;
   static ctype generate_root(int len, bool invert) {
       ctype wlen = invert ? inv_root : root;
       for (int i = len; i < root_pw; i <<= 1)</pre>
           wlen = (111 * wlen * wlen) % mod;
       return wlen;
   static void prepare_output(
       const cvector& vin, ivector& vout, uint32_t out_size) {
       vout = vin;
```

```
while (vout.size() > 1 && vout.back() == 0) vout.pop_back();
}
static void invert(cvector& a) {
    int32_t inv_n = inverse(a.size(), mod);
    for (auto& x : a) x = (111 * x * inv_n) % mod;
}
static int32_t root, inv_root, root_pw, mod;
};
// Let mod = c * 2^NBIT + 1. Then, NTT::root is
// (g^c) % mod, where g is primitive root of mod.
int32_t NTT::root = /* ... */
int32_t NTT::inv_root = inverse(NTT::root, modP);
int32_t NTT::mod = modP;
```

#### 7.4 FFT

```
const double PI = 2 * acos(0);
struct C {
 double a, b;
 C(): a(0), b(0) {}
 C (double a, double b) : a(a), b(b) {}
 C (double theta) : a(cos(theta)), b(sin(theta)) {}
 C bar () const { return C(a, -b); }
 double modsq () const { return a * a + b * b; }
 C operator+ (const C &c) const { return C(a + c.a, b + c.b); }
 C operator* (const C &c) const { return C(a * c.a - b * c.b, a * c.b
      + b * c.a); }
 C operator/ (const C &c) const {
   C r = (*this) * c.bar();
   return C(r.a / c.modsq(), r.b / c.modsq());
 }
};
// O(nlogn)
// dir is direction of Fourier transform
void fft (C *in, C *out, int step, int size, int dir) {
 if (size < 1) return;
 if (size == 1) { out[0] = in[0]; return; }
 fft(in, out, step*2, size/2, dir);
 fft(in + step, out + size/2, step*2, size/2, dir);
 for (int i = 0; i < size/2; i++) {
   C even = out[i], odd = out[i + size/2];
```

```
out[i] = even + C(dir * 2*PI * i / size) * odd;
  out[i + size/2] = even + C(dir * 2*PI * (i + size/2) / size) * odd;
}

// c[i] = sum of a[j] * b[i-j]

// n is power of 2; index is cyclic
void convolve (int n, C *a, C *b, C *c) {
  C *fa = new C[n];
  C *fb = new C[n];
  C *fc = new C[n];
  fft(a, fa, 1, n, 1);
  fft(b, fb, 1, n, 1);
  for (int i = 0; i < n; i++) fc[i] = fa[i] * fb[i];
  fft(fc, c, 1, n, -1);
  for (int i = 0; i < n; i++) c[i] = c[i] / C(n,0);
}</pre>
```

## 8 Geometry

```
double EPS = 1e-8;
double PI = acos(-1.0);

bool equal (double x, double y) { return fabs(x - y) < EPS; }
int sign (double x) {
  if (equal(x, 0.0)) return 0;
  return x > 0.0 ? 1 : -1;
}
```

## 8.1 Point

```
struct Point {
   double x, y;

Point (double x, double y) : x(x), y(y) {}

friend bool operator== (Point p, Point q) { return equal(p.x, q.x)
        && equal(p.y, q.y); }

friend Point operator+ (Point p, Point q) { return Point(p.x + q.x,
        p.y + q.y); }

friend Point operator- (Point p, Point q) { return Point(p.x - q.x,
        p.y - q.y); }
```

```
friend Point operator* (Point p, double k) { return Point(p.x * k,
    p.y * k); }
friend Point operator/ (Point p, double k) { return p * (1.0 / k); }
static double arg (Point p) { return atan2(p.y, p.x); }
static double norm (Point p) { return sqrt(p.x * p.x + p.y * p.y); }
static double dot (Point p, Point q) { return p.x * q.x + p.y * q.y;
    }
static double cross (Point p, Point q) { return p.x * q.y - q.x *
static double dist (Point p, Point q) { return norm(p - q); }
static double det (Point p, Point q, Point r) { return cross(q-p,
static Point rotate (Point p, double theta) {
 return Point(p.x * cos(theta) - p.y * sin(theta), p.x * sin(theta)
      + p.y * cos(theta));
/* triangle */
static Point mass_center (Point p1, Point p2, Point p3) {
 return (p1 + p2 + p3) / 3.0;
static Point outer_center (Point p1, Point p2, Point p3) {
 double a1 = p2.x - p1.x, b1 = p2.y - p1.y, c1 = (a1*a1+b1*b1) /
 double a2 = p3.x - p1.x, b2 = p3.y - p1.y, c2 = (a2*a2+b2*b2) /
      2.0:
  double d = a1 * b2 - a2 * b1;
 double x = p1.x + (c1*b2 - c2*b1) / d;
 double y = p1.y + (a1*c2 - a2*c1) / d;
 return Point(x, y);
static Point outer_center (Point p1, Point p2) {
 return (p1 + p2) / 2.0;
static Point ortho_center (Point p1, Point p2, Point p3) {
 return mass_center(p1, p2, p3) * 3.0 - outer_center(p1, p2, p3) *
      2.0:
static Point inner_center (Point p1, Point p2, Point p3) {
 double a = dist(p2, p3);
 double b = dist(p3, p1);
 double c = dist(p1, p2);
 return (p1 * a + p2 * b + p3 * c) / (a + b + c);
/* triangle */
```

```
// divide and conquer: O(nlogn)
// tested on HDU 1007
static pair <double, pair <Point, Point >> closest_pair (vector <Point >>
    ps) {
 int n = ps.size();
 vector<int> rank(n);
 for (int i = 0; i < n; i++) rank[i] = i;
 sort(rank.begin(), rank.end(), [&ps](int i, int j) -> bool {
      return ps[i].x < ps[j].x; });
 return closest_pair(ps, rank, 0, n);
static pair<double,pair<Point,Point>> closest_pair (vector<Point>
    &ps, vector<int> &rank, int 1, int r) {
 auto ans_cmp = [](pair<double,pair<Point,Point>> i,
      pair<double,pair<Point,Point>> j) -> bool { return i.first <</pre>
     j.first; };
 if (r - 1 < 20) {
   pair<double,pair<Point,Point>> ans = {0x7ffffffff, {Point(0,0),
        Point(0,0)}};
   for (int i = 1; i < r; i++) {
     for (int j = i+1; j < r; j++) {
       if (ans.first > dist(ps[rank[i]], ps[rank[j]])) {
         ans = {dist(ps[rank[i]], ps[rank[j]]), {ps[rank[i]],
             ps[rank[j]]}};
       }
     }
   }
   return ans;
 int mid = (1 + r) / 2;
 auto ans = min(closest_pair(ps, rank, 1, mid), closest_pair(ps,
      rank, mid, r), ans_cmp);
 int tl; for (tl = 1; ps[rank[tl]].x < ps[rank[mid]].x - ans.first;</pre>
 int tr; for (tr = r-1; ps[rank[tr]].x > ps[rank[mid]].x +
      ans.first; tr--);
 sort(rank.begin()+tl, rank.begin()+tr, [&ps](int i, int j) -> bool
      { return ps[i].y < ps[j].y; });
 for (int i = tl; i < tr; i++) {
   for (int j = i+1; j < min(tr, i+6); j++) {
     if (ans.first > dist(ps[rank[i]], ps[rank[j]])) {
       ans = {dist(ps[rank[i]], ps[rank[j]]), {ps[rank[i]],
           ps[rank[j]]}};
     }
   }
```

```
}
  sort(rank.begin()+tl, rank.begin()+tr, [&ps](int i, int j) -> bool
      { return ps[i].x < ps[j].x; });
 return ans;
}
// farthest pair in a convex hull
// DEBUG: maybe not good at when all points are colinear
// tested on POJ 2187
static pair<double,pair<Point,Point>> farthest_pair (vector<Point>)
  auto ans_cmp = [](pair<double,pair<Point,Point>> i,
      pair<double,pair<Point,Point>> j) -> bool { return i.first <</pre>
      j.first; };
 int n = ps.size();
 pair<double,pair<Point,Point>> ans = {0.0, {Point(0,0),
      Point(0,0)}};
 if (n == 1) return ans;
 for (int i = 0, j = 1; i < n; i++) {
   while (sign(det(ps[i], ps[(i+1)%n], ps[j]) - det(ps[i],
        ps[(i+1)\%n], ps[(j+1)\%n])) == -1) {
     j = (j+1)%n;
   }
   ans = max(ans, {dist(ps[i], ps[j]), {ps[i], ps[j]}}, ans_cmp);
   ans = \max(ans, \{dist(ps[(i+1)\%n], ps[(j+1)\%n]), \{ps[(i+1)\%n], \}
        ps[(j+1)%n]}}, ans_cmp);
 }
 return ans;
// Graham scan: O(nlogn); result in counter-clockwise
// tested on POJ 2187 indirectly
static vector<Point> convex_hull (vector<Point> ps) {
 int n = ps.size();
 if (n < 3) return ps;
 for (int i = 1; i < n; i++) {
   if (ps[0].y > ps[i].y \mid | (ps[0].y == ps[i].y && ps[0].x >
        ps[i].x)) {
     swap(ps[0], ps[i]);
   }
 }
 Point base = ps[0];
  sort(ps.begin()+1, ps.end(), [&](Point p, Point q) -> bool {
      return det(base, p, q) > 0 || (det(base, p, q) == 0 &&
      dist(base, p) < dist(base, q)); });</pre>
  vector<Point> ans = \{ps[0], ps[1], ps[2]\};
```

```
for (int i = 3; i < n; i++) {
    while (sign(det(ans[ans.size()-1], ans[ans.size()-2], ps[i])) ==
        1) ans.pop_back();
    ans.push_back(ps[i]);
    }
    return ans;
}
</pre>
```

#### 8.2 Line

```
struct Line {
 Point a, b;
 Line (Point a, Point b) : a(a), b(b) {}
 static double dist (Line 1, Point p) {
   return fabs(Point::det(p, l.a, l.b) / Point::dist(l.a, l.b));
 static Point proj (Line 1, Point p) {
   double r = Point::dot(1.b - 1.a, p - 1.a) / Point::dot(1.b - 1.a,
       1.b - 1.a);
   return 1.a * (1 - r) + 1.b * r:
 }
 static bool on_segment (Line 1, Point p) {
   return sign(Point::det(p, 1.a, 1.b)) == 0 && sign(Point::dot(p -
       1.a, p - 1.b)) <= 0;
 static bool parallel (Line 1, Line m) {
   return sign(Point::cross(l.a - l.b, m.a - m.b)) == 0;
 }
 static Point line_x_line (Line 1, Line m) {
   double s1 = Point::det(m.a, l.a, m.b);
   double s2 = Point::det(m.a, 1.b, m.b);
   return (1.b * s1 - 1.a * s2) / (s1 - s2);
 static bool two_segments_intersect (Line 1, Line m) {
   double dla = Point::det(1.b, m.a, m.b);
   double dlb = Point::det(l.a, m.a, m.b);
```

```
double dma = Point::det(m.b, l.a, l.b);
 double dmb = Point::det(m.a, l.a, l.b);
 if (sign(dla * dlb) == -1 && sign(dma * dmb) == -1) return true;
 if (sign(dla) == 0 && on_segment(m, 1.b)) return true;
 if (sign(dlb) == 0 && on_segment(m, 1.a)) return true;
 if (sign(dma) == 0 && on_segment(1, m.b)) return true;
 if (sign(dmb) == 0 && on_segment(1, m.a)) return true;
 return false;
static bool any_segments_intersect (vector<Line> ls) {
 vector<pair<Point,pair<int,int>>> items;
 for (int i = 0; i < ls.size(); i++) {
   Line &l = ls[i];
   if (1.a.x > 1.b.x) swap(1.a, 1.b);
   items.push_back({1.a, {0, i}});
   items.push_back({1.b, {1, i}});
 }
 sort(items.begin(), items.end(), [](pair<Point,pair<int,int>> a,
      pair<Point,pair<int,int>> b) -> bool {
   if (sign(a.first.x - b.first.x) == -1) return true;
   if (sign(a.first.x - b.first.x) == 1) return false;
   if (a.second.first < b.second.first) return true;</pre>
   if (a.second.first > b.second.first) return false;
   return a.first.y < b.first.y;</pre>
 });
 auto cmp = [&](int i, int j) -> bool { return ls[i].a.y <
      ls[j].a.y; };
 set<int,decltype(cmp)> s(cmp);
 for (auto &item : items) {
   if (item.second.first == 0) {
     auto it = s.insert(item.second.second).first;
     int id = *it:
     int prev_id = (it == s.begin()) ? -1 : *(prev(it));
     int next_id = (next(it) == s.end()) ? -1 : *(next(it));
     if (prev_id != -1 && two_segments_intersect(ls[id],
         ls[prev_id])) return true;
     if (next_id != -1 && two_segments_intersect(ls[id],
         ls[next_id])) return true;
   } else {
     auto it = s.find(item.second.second);
     int id = *it;
     int prev_id = (it == s.begin()) ? -1 : *(prev(it));
     int next_id = (next(it) == s.end()) ? -1 : *(next(it));
     if (prev_id != -1 && next_id != -1 &&
         two_segments_intersect(ls[prev_id], ls[next_id])) return
```

```
true;
    s.erase(it);
}
return false;
}
;
```

## 8.3 Halfplane

```
struct HalfPlane {
 Point s, t; // half plane on the left of ray from p to q
 HalfPlane (Point s, Point t) : s(s), t(t) {}
 double eval (Point p) {
   double a, b, c; // ax+by+c \le 0
   a = t.y - s.y;
   b = s.x - t.x;
   c = Point::cross(t, s);
   return p.x * a + p.y * b + c;
 static Point halfplane_x_line (HalfPlane hp, Line 1) {
   Point p = 1.a, q = 1.b:
   double vp = hp.eval(p), vq = hp.eval(q);
   double x = (vq * p.x - vp * q.x) / (vq - vp);
   double y = (vq * p.y - vp * q.y) / (vq - vp);
   return Point(x, y);
 }
 static vector<Point> halfplanes_x (vector<HalfPlane> hps) {
   sort(hps.begin(), hps.end(), [](HalfPlane a, HalfPlane b) -> bool {
    int sgn = sign(Point::arg(a.t - a.s) - Point::arg(b.t - b.s));
    return sgn == 0 ? (sign(b.eval(a.s)) == -1) : (sgn < 0);
   });
   deque<HalfPlane> q {hps[0]};
   deque<Point> ans;
   for (int i = 1; i < hps.size(); i++) {
    if (sign(Point::arg(hps[i].t - hps[i].s) - Point::arg(hps[i-1].t
         - hps[i-1].s) == 0) continue;
    while (ans.size() > 0 && sign(hps[i].eval(ans.back())) == 1) {
         ans.pop_back(); q.pop_back(); }
    while (ans.size() > 0 && sign(hps[i].eval(ans.front())) == 1) {
         ans.pop_front(); q.pop_front(); }
```

## 8.4 Polygon

```
struct Polygon {
 int n;
 vector<Point> p; // always counter-clockwise
 Polygon (vector<Point> p) : p(p), n(p.size()) {}
 double perimeter () {
   double ans = 0;
   for (int i = 0; i < n; i++) {
     ans += Point::dist(p[i], p[(i+1)%n]);
   }
   return ans;
 double area () {
   double ans = 0;
   for (int i = 1; i < n-1; i++) {
     ans += Point::det(p[0], p[i], p[i+1]) / 2.0;
   return ans;
 Point mass_center () {
   Point ans(0.0, 0.0);
   double a = area();
   if (sign(a) == 0) return ans;
   for (int i = 1; i < n-1; i++) {
```

```
ans = ans + ((p[0] + p[i] + p[i+1]) / 3.0) * (Point::det(p[0],
        p[i], p[i+1]) / 2.0);
 return ans / a:
}
// first is grid point inside polygon; second is grid point on edge.
    vertices has to be grid points
pair<int,int> grid_point_cnt () {
 int first = 0, second = 0;
  for (int i = 0; i < n; i++) {
   second += gcd(abs((int)(p[(i+1)%n].x - p[i].x)),
        abs((int)(p[(i+1)%n].y - p[i].y)));
  first = (int)area() + 1 - second / 2;
  return {first, second};
}
int gcd(int p, int q) \{ return q == 0 ? p : gcd(q, p%q); \}
bool is_simple_convex_polygon () {
 for (int i = 0; i < n; i++) { // convexity
   if (sign(Point::det(p[i], p[(i+1)\%n], p[(i+2)\%n])) == -1) return
        false;
 }
  for (int i = 1; i < n-1; i++) { // simplicity
   if (sign(Point::det(p[0], p[i], p[i+1])) == -1) return false;
 }
 return true;
}
// O(n)
// returns 1 for in, 0 for on, -1 for out
static int point_in_polygon (Polygon po, Point p0) {
 int cnt = 0;
 for (int i = 0; i < po.n; i++) {
   if (Line::on_segment(Line(po.p[i], po.p[(i+1)%po.n]), p0)) return
   int k = sign(Point::det(p0, po.p[i], po.p[(i+1)%po.n]));
   int d1 = sign(po.p[i].y - po.y);
   int d2 = sign(po.p[(i+1)\%po.n].y - po.y);
   if (k == 1 && d1 != 1 && d2 == 1) cnt++;
   if (k == -1 && d2 != 1 && d1 == 1) cnt--;
  return cnt ? 1 : -1;
}
```

```
// O(log(n))
// returns 1 for in, 0 for on, -1 for out
static int point_in_convex_polygon (Polygon po, Point p0) {
 Point point = (po.p[0] + po.p[po.n/3] + po.p[2*po.n/3]) / 3.0;
 int l = 0, r = po.n;
 while (r - 1 > 1) {
   int mid = (1 + r) / 2;
   if (sign(Point::det(point, po.p[1], po.p[mid])) == 1) {
     if (sign(Point::det(point, po.p[1], p0)) != -1 &&
          sign(Point::det(point, po.p[mid], p0)) == -1) r = mid;
     else 1 = mid;
   } else {
     if (sign(Point::det(point, po.p[1], p0)) == -1 &&
         sign(Point::det(point, po.p[mid], p0)) != -1) l = mid;
     else r = mid:
   }
 }
 r %= po.n;
 return -sign(Point::det(p0, po.p[r], po.p[l]));
Polygon convex_polygon_x_halfplane (HalfPlane hp, Polygon po) {
 vector<Point> ps;
 for (int i = 0; i < po.n; i++) {
   if (sign(hp.eval(po.p[i])) == -1) {
     ps.push_back(po.p[i]);
   } else {
     if (sign(hp.eval(po.p[(i-1+po.n)\%po.n])) == -1) {
       ps.push_back(HalfPlane::halfplane_x_line(hp, Line(po.p[i],
           po.p[(i-1+po.n)%po.n])));
     if (sign(hp.eval(po.p[(i+1)\%po.n])) == -1) {
       ps.push_back(HalfPlane::halfplane_x_line(hp, Line(po.p[i],
           po.p[(i+1)%po.n])));
     }
   }
 }
 return Polygon(ps);
static Polygon convex_polygon_x_convex_polygon (Polygon po1, Polygon
    po2) {
 vector<HalfPlane> hps;
 for (int i = 0; i < po1.n; i++) {
   hps.push_back(HalfPlane(po1.p[i], po1.p[(i+1)%po1.n]));
 }
```

```
for (int i = 0; i < po2.n; i++) {
   hps.push_back(HalfPlane(po2.p[i], po2.p[(i+1)%po2.n]));
}
return Polygon(HalfPlane::halfplanes_x(hps));
}
};</pre>
```

#### 8.5 Circle

```
struct Circle {
  Point center;
 double radius;
 Circle (Point center, double radius) : center(center),
      radius(radius) {}
  static bool in_circle (Circle c, Point p) {
   return sign(Point::dist(p, c.center) - c.radius) == -1;
  static Circle min_circle_cover (vector<Point> p) {
   Circle ans(p[0], 0.0);
   random_shuffle(p.begin(), p.end());
   for (int i = 1; i < p.size(); i++) if (!in_circle(ans, p[i])) {</pre>
     ans.center = p[i]; ans.radius = 0;
     for (int j = 0; j < i; j++) if (!in_circle(ans, p[j])) {
       ans.center = Point::outer_center(p[i], p[j]);
       ans.radius = Point::dist(p[j], ans.center);
       for (int k = 0; k < j; k++) if (!in_circle(ans, p[k])) {
         ans.center = Point::outer_center(p[i], p[j], p[k]);
         ans.radius = Point::dist(p[k], ans.center);
     }
   }
   return ans;
 }
};
```

## 8.6 Simplex volume

```
// AB AC AD BC BD CD
double simplex_volume (double 1, double n, double a, double m, double b, double c) {
```

## 8.7 Count gridpoints under a line

```
// Count gridpoints under a line
// Compute for (int i = 0; i < n; i++) s += floor((a+b*i)/m);
long long count_gridpoints (long long n, long long a, long long b,
        long long m) {
   if (b == 0) return n * (a / m);
   if (a >= m) return n * (a / m) + count_gridpoints(n, a%m, b, m);
   if (b >= m) return (n-1) * n / 2 * (b / m) + count_gridpoints(n, a,
        b%m, m);
   return count_gridpoints((a+b*n)/m, (a+b*n)%m, m, b);
}
```

## 8.8 Simpson's Union Of Circles

```
struct dot
   double x, y;
   double dis(dot &o)
       return sqrt(sqr(x - o.x) + sqr(y - o.y));
};
int 1x = 1000, rx = -1000;
struct circle
{
   dot o; int r;
   void init()
   {
       int x, y;
       scanf("%d%d%d", &x, &y, &r);
       lx = min(lx, x - r); rx = max(rx, x + r);
       o.x = x; o.y = y;
   }
   bool in(circle &b)
   {
```

```
return (b.r - r - o.dis(b.o) \ge -eps);
   }
    bool operator==(const circle &b)
       return r == b.r \&\& fabs(o.x - b.o.x) \le eps \&\& fabs(o.y - b.o.x)
            b.o.y) <= eps;
}tmp[Maxn], c[Maxn];
struct seg
{
    double v; int s;
    bool operator<(const seg &o)</pre>
       const{return v < o.v - eps;}</pre>
1[Maxn * 2];
int n, m;
void Init()
    scanf("%d", &m);
    for (int i = 1; i \le m; ++i)
       tmp[++n].init();
       for (int j = 1; j \le n - 1; ++j)
           if (tmp[j] == tmp[n])
               {--n; break;}
    m = n; n = 0;
    for (int i = 1; i \le m; ++i)
       bool f = 0;
       for (int j = 1; j \le m; ++j) if (j != i)
           if (tmp[i].in(tmp[j]))
           {
               f = 1;
               break;
       if (!f) c[++n] = tmp[i];
   }
}
inline double get(double x)
    int t = 0, now = 0;
    double d, last, s = 0;
    for (int i = 1; i \le n; ++i)
    {
```

```
if (fabs(x - c[i].o.x) - c[i].r \ge -eps) continue;
       d = sqrt(sqr(c[i].r) - sqr(x - c[i].o.x));
       l[++t].v = c[i].o.v - d; l[t].s = 1;
       l[++t].v = c[i].o.v + d; l[t].s = -1;
   }
   sort(1 + 1, 1 + 1 + t);
   for (int i = 1; i \le t; ++i)
       now += l[i].s:
       if (now == 1 && l[i].s == 1) last = l[i].v;
       if (now == 0) s += 1[i].v - last;
   }
   return s;
}
double simpson(double 1, double r, double 1x, double mx, double rx)
   double m = (1 + r) * 0.5, lp, rp, s, ls, rs;
   lp = get((1 + m) * 0.5);
   rp = get((m + r) * 0.5);
   s = (1x + rx + 4 * mx) * (r - 1) / 6;
   ls = (lx + mx + 4 * lp) * (m - 1) / 6;
   rs = (mx + rx + 4 * rp) * (r - m) / 6;
   if (fabs(ls + rs - s) \le 1e-6)
   return simpson(1, m, lx, lp, mx) + simpson(m, r, mx, rp, rx);
}
void Work()
   double s = 0, last = get(lx), now;
   for (int i = lx; i \le rx - 1; ++i)
       now = get(i + 1);
       if (fabs(last) > eps || fabs(now) > eps)
           s += simpson(i, i + 1, last, get(i + 0.5), now);
       last = now;
   printf("%.3lf\n", s);
}
```

## 9 Misc

#### 9.1 Date

#### 9.1.1 Date to Day of Week

#### 9.1.2 Count Days from AD

```
const int days = 365;
const int s[] = {0, 31, 28, 31, 30, 31, 30, 31, 30, 31, 30, 31};
bool IsLeap (int y) {
  return y % 400 == 0 || (y % 100 && y % 4 == 0);
}
int leap (int y) {
  return y/4 - y/100 + y/400;
}
int calc (int day, int mon, int year) {
  int res = (year-1) * days + leap(year-1);
  for (int i = 1; i < mon; ++i) res += s[i];
  if (IsLeap(year) && mon > 2) res++;
  res += day;
  return res;
}
```

## ICPC Math Table

\* p is prime

## 1 Number Theory

#### Fermat's little theorem

$$a^{p-1} \equiv 1 \pmod{p}$$
  
 $a^{\phi(n)} \equiv 1 \pmod{n}$  where  $\gcd(a, n) = 1$   
 $a^m \equiv a^{m\%\phi(n)+\phi(n)} \pmod{n}$ 

### Euler's totient function

$$\overline{\phi(n)} = |\{x \mid 1 \le x \le n, \gcd(x, n) = 1\}|$$

$$\phi(n) = n \prod_{p \mid n} (1 - \frac{1}{p})$$

$$\phi(mn) = \phi(m)\phi(n) \text{ if } \gcd(m, n) = 1$$

$$\phi(mn) = \phi(m)\phi(n) \frac{d}{\phi(d)} \text{ where } d = \gcd(m, n)$$

$$\phi(m)\phi(n) = \phi(lcm(m, n))\phi(\gcd(m, n))$$

$$\sum_{d \mid n} \phi(d) = n$$

$$\sum_{d \mid n} \frac{n}{d}\phi(d) = \sum_{k=1..n} \gcd(k, n)$$

$$\phi(n)d(n) = \sum_{k=1..n} \gcd(k, n)$$

$$\phi(n)d(n) = \sum_{k=1..n} \gcd(k, n)$$

$$\frac{1}{2}n\phi(n) = \sum_{k=1..n} k$$

$$a \mid b \to \phi(a) \mid \phi(b)$$

$$n \mid \phi(a^n - 1) \text{ for } a, n > 1$$

### Mobius function

$$\mu(n) = \begin{cases} 0 \text{ if } n \text{ has squared prime factor} \\ 1 \text{ if } n \text{ has even } \# \text{ of prime factors} \\ -1 \text{ if } n \text{ has odd } \# \text{ of prime factors} \end{cases}$$

$$\begin{split} \sum_{d|n} \mu(d) &= [n == 1] \\ n \sum_{d|n} \frac{\mu(d)}{d} &= \phi(n) \\ \sum_{d|n} \frac{\mu^2(d)}{\phi(d)} &= \frac{n}{\phi(n)} \\ \forall n, g(n) &= \sum_{d|n} f(d) \rightarrow \forall n, f(n) = \sum_{d|n} \mu(d) g(\frac{n}{d}) \end{split}$$

## **Primality criteria** (p is prime iff)

$$\prod_{1 \le k \le p-1} (2^k - 1) \equiv p \mod (2^p - 1)$$
$$(p-1)! \equiv -1 \mod p$$

## 2 Combinatorics

$$\binom{n}{0} + \dots + \binom{n}{n} = 2^{n}$$

$$\binom{n}{0} + \binom{n}{2} + \dots = 2^{n-1}$$

$$\binom{n}{1} + \binom{n}{3} + \dots = 2^{n-1}$$

$$0\binom{n}{0} + \dots + n\binom{n}{n} = n2^{n-1}$$

$$0^{2}\binom{n}{0} + \dots + n^{2}\binom{n}{n} = n(n+1)2^{n-2}$$

$$n\binom{n-1}{k-1} = k\binom{n}{k}$$

$$\binom{n-1}{k} + \binom{n-1}{k-1} = \binom{n}{k}$$

$$\binom{k}{k} + \dots + \binom{n}{k} = \binom{n+1}{k+1}$$

$$\binom{m}{0}\binom{n}{k} + \dots + \binom{m}{k}\binom{n}{0} = \binom{m+n}{k}$$

$$\binom{n}{0}^{2} + \dots + \binom{n}{n}^{2} = \binom{2n}{n}$$

$$\text{Lucas: } \binom{m}{n} \equiv \prod \binom{m_{i}}{n_{i}} \pmod{p}$$
Wolstenholme: 
$$\binom{2p-1}{p-1} \equiv 1 \pmod{p^{3}} \text{ where } p > 3$$
Wolstenholme: 
$$\binom{ap}{bp} \equiv \binom{a}{b} \pmod{p^{3}} \text{ where } p > 3$$

```
# lower-diagonal paths from (0,0) to (n,m) (n \ge m) = \frac{n-m+1}{n+1} \binom{n+m}{m}
Lex-order index (1-based) of r-subset \{a_1..a_r\} of \{1..n\} = \binom{n}{r} - \binom{n-a_1}{r} - \dots - \binom{n-a_r}{1}
```

Enum r-subsets of n-set in lex-order

```
int a[] = {1...r}
while (1) {
    int k;
    for (k = r; k > 0 && !(a[k] < n && a[k]+1 != a[k]); k--);
    if (k == 0) break;
    for (int i = r; i >= k; i--) a[i] = a[k] + (i - k + 1);
}
```

Enum r-subsets of n-set

**Difference table** leftmost diagonal =  $c_0, \ldots c_p, 0, \ldots \rightarrow$  original sequence

$$h_n = c_0 \binom{n}{0} + \dots + c_p \binom{n}{p}$$
$$\sum_{k=0\dots n} h_k = c_0 \binom{n+1}{1} + \dots + c_p \binom{n+1}{p+1}$$

### Catalan number

 $\overline{C_n = \# \pm 1}$  sequences with non-negative prefix sum

$$C_n = \frac{1}{n+1} \binom{2n}{n}$$
$$C_n = \frac{4n-2}{n+1} C_{n-1}$$

### Stirling-1 number

 $\overline{s(p,k)} = \# p \text{ diff items}$  into k same circular permutations

$$s(p,0) = 0 (p \ge 1)$$

$$s(p,p) = 1 (p \ge 0)$$

$$s(p,k) = (p-1)s(p-1,k) + s(p-1,k-1) (1 \le k \le p-1)$$

$$A_n^p = \sum_{k=0..p} (-1)^{p-k} s(p,k) n^k$$

## Stirling-2 number

 $\overline{S(p,k)} = \# p$  diff items into k same boxes, no empty box

$$S(p,0) = 0 (p \ge 1)$$

$$S(p,p) = 1 (p \ge 0)$$

$$S(p,k) = kS(p-1,k) + S(p-1,k-1) (1 \le k \le p-1)$$

$$S(p,k) = \frac{1}{k!} \sum_{i=0..k} (-1)^i \binom{k}{i} (k-i)^p$$

$$n^p = \sum_{k=0..p} S(p,k) A_n^k$$

# p diff items into k diff boxes = k!S(p,k)

## $\mathbf{Bell}\ \mathbf{number}$

 $\overline{B_p = \# p \text{ diff items into same boxes}}$ 

$$B_p = S(p,0) + \dots + S(p,p)$$

$$B_p = {p-1 \choose 0} B_0 + \dots + {p-1 \choose p-1} B_{p-1}$$

$$B_{p^i+k} \equiv iB_k + B_{k+1} \pmod{p}$$

## Generating function

r-combination:  $\prod (1+x^1+x^2+\ldots+x^{f_i})$ r-arrangement:  $r! \prod (1+\frac{x^1}{1!}+\frac{x^2}{2!}+\ldots+\frac{x^{f_i}}{f_i!})$ 

Integer partition:  $\prod_{k=1..n} (1-x^k)^{-1}$ 

## Burnside lemma, Polya enum theorem

# inequivalent colorings on n-set under a permutation group.

$$N(C,G) = \frac{1}{|G|} \sum_{f \in G} |C(f)| = \frac{1}{|G|} \sum_{f \in G} k^{\#(f)} = \frac{1}{|G|} \sum_{f \in G} k^{\sum e_i}$$

G is the equivalent permutation group

C is all colorings on n-set

N(C,G) is # inequivalent colorings

C(f) is the stable kernel of permutation f

k is the number of colors available

#(f) is the number of cycles in permutation f

 $e_1 \dots e_n$  is the type of permutation f - it has  $e_i$  i-cycles

## 3 Graph Theory

## Havel-Hakimi algo

degree sequence  $(d_1 \ge ... \ge d_n)$  is simple-graphic iff  $(d_2-1...d_{d_1+1}-1, d_{d_1+2}...d_n)$  is simple-graphic. Equivalently, connect largest-degree node with other largest-degree nodes.

Erdos-Gallai theorem:  $(d_1 \ge ... \ge d_n)$  is simple-graphic iff

$$\forall k \in [1, n] \sum_{i=1}^{k} d_i \le k(k-1) + \sum_{i=k+1}^{n} \min(d_i, k)$$

## $\label{eq:Vizing} \mbox{Vizing's theorem} + \mbox{Misra-Gries edge coloring algo}$

adjacent edges cannot have same color, uses  $\max(deq(v)) + 1$  colors.

## 4 Game Theory

Nim Lose iff XOR sum is zero

### SG function

P-position: first lose

N-position: second lose

Final node must be P

N's successors contain at least one P

P's successors contain all N

 $SG(x) = mex(\{SG(y) \mid y \text{ is successor of } x \})$ 

SG(x) = 0 iff x is P-position

Composite game's SG value is the XOR sum of simple games

## 5 Numerical Methods

**Newton's method** solve f(x) = 0 by  $x \leftarrow x - f(x)/f'(x)$ 

## 6 Miscellaneous

$$\gcd(2^a - 1, 2^b - 1) = 2^{\gcd(a,b)} - 1$$

 $x^2 + y^2 = n$  has integer solution  $\leftrightarrow n = \prod p_i^{e_i}$ , there are no i s.t.  $p_i \equiv 3 \pmod{4}$  and  $e_i \equiv 1 \pmod{2}$ 

## Fibbonacci

$$\gcd(F_n, F_m) = F_{\gcd(n,m)}$$
$$b \mid a \leftrightarrow F_b \mid F_a$$

## Derangements

$$D_n = n!(1 - \frac{1}{1!} + \frac{1}{2!} - \dots + \frac{(-1)^n}{n!})$$

$$D_n = (n-1)(D_{n-2} + D_{n-1})$$

$$D_n = nD_{n-1} + (-1)^n$$

Gray sequence 
$$G_i = i \operatorname{xor} (i >> 1)$$

Farey sequence sorted  $\frac{a}{b}$   $(1 \le a < b \le N, \gcd(a, b) = 1)$ 

$$\begin{aligned} \frac{a_0}{b_0} &= \frac{0}{1} \\ \frac{a_1}{b_1} &= \frac{1}{N} \\ \frac{a_n}{b_n} &= \frac{a_{n-1} \lfloor \frac{N+b_{n-2}}{b_{n-1}} \rfloor - a_{n-2}}{b_{n-1} \lfloor \frac{N+b_{n-2}}{b_{n-1}} \rfloor - b_{n-2}} \end{aligned}$$

**Dilworth theorem** fewest chain split = longest reverse chain