
Redes de Computadores

Introduction

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Introduction to the Course

RCOM – Professors, Language

- ◆ Professors
 - » Manuel Ricardo
 - » Pedro Brandão
 - » Maria Teresa Andrade
 - » Rui Lopes Campos
 - » Sérgio Lopes Crisóstomo
 - » Filipe Borges Teixeira
- ◆ Information about RCOM
 - available in moodle
- ◆ Language
 - » Slides and books in English
 - » Lectures in Portuguese
 - » *Suitable for English-speaking students*

Bibliografia

- ◆ Main book
 - Andrew Tanenbaum, David Wetherall,
Computer Networks, 5/E
Prentice Hall
2011**
- ◆ Slides presented in classes
 - » Follow the main book
 - » Complemented with information from other sources
 - » Oriented to fundamentals; details in book

Bibliografia – Other books

- ◆ **Dimitri Bertsekas, Robert Gallager, Data Networks, 2nd Edition, 1992, Prentice Hall**
 - » Oriented to the fundamental aspects of data networks with formal (math) descriptions
 - » Available also in <http://web.mit.edu/dimitrib/www/datanets.html>
 - » Examples on outdated networks
- ◆ **Larry L. Peterson, Bruce S. Davie, Computer Networks - A Systems Approach, 4th Edition, 2007, Morgan Kaufmann**
 - » Less generic than Tanenbaum; oriented to TCP/IP and implementation aspects
- ◆ **James F. Kurose, Keith W. Ross, Computer Networking - a Top-Down Approach, 2010, 5th Edition, Pearson**
 - Similar to Tanenbaum; uses top-down approach; more focused on applications than in physical layer
- ◆ **W. Richard Stevens, TCP/IP Illustrated: The Protocols (Vol. 1), 1994, Addison-Wesley.**
 - » The book of TCP/IP stack
- ◆ **William Stallings, Data & Computer Communications, 8th Edition, 2007, Prentice Hall**
 - » Generic and good book; addresses also telecom networks

Types of Classes

- ◆ *Aulas teóricas*

- » Oriented to the fundamental aspects of Computer Networks
- » Additional **reading required at home**
- » **Weekly homeworks**

questions to be answered before next lecture **through moodle**

- ◆ *Aulas laboratoriais*

2 laboratory projects

- » 1st lab: protocol development, Linux, C programming, file transfer
- » 2nd lab: configuration computer network (switches, routers, computers)

Evaluation of RCOM

Frequênciа

- ◆ L1 - grade of 1st lab
- ◆ L2 - grade of 2nd lab
- ◆ H - grade of homeworks
- ◆ FQ - grade of FREQUÊNCIA

- ◆ $FQ = 0,4*L1 + 0,4*L2 + 0,2*H$

- ◆ if ($FQ < 8,0$) FQ = "No Admission to Exams"

Classificaçāo Final

- ◆ E - grade of final exam
- ◆ FQ - grade of FREQUÊNCIA
- ◆ AD - grade of distributed evaluation
- ◆ CF - final grade

- ◆ if ($FQ > E + 5$) $AD = E + 5$ else $AD = FQ$

- ◆ $CF = 0.4 * AD + 0.6 * E$

- ◆ if ($E < 8.0$) $CF = E$

Learning objectives

- ◆ Fundaments of network design and analysis
 - » Communication channels and data link control
 - » Delay and loss models in data networks
 - » Multi-access communications
 - » Routing in computer networks
 - » Flow and congestion control
- ◆ Technologies in use
 - » Ethernet, WLAN, Internet, TCP/IP communications stack
- ◆ Implementation
 - » Protocol development in UNIX
 - » Network configuration

Introduction to Computer Networks

-
- » *What are the **main uses** of computer networks?*
 - » *What are the main **types of networks**?*
 - » *What is a **protocol**? What is a **service**?*
 - » *What is a **protocol stack**?*
 - » *What are the communication layers of the **Internet reference model**?*
 - » *What are the differences between **circuit switching** and **packet switching**?*
 - » *What is the **propagation delay**, T_{prop} ?*
 - » *What is the **packet transmission delay**, T_{pac} ?*

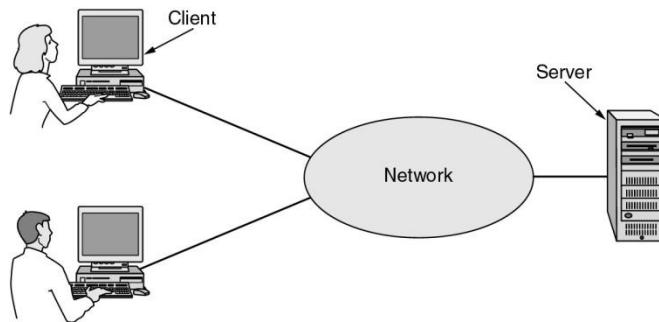
Uses of Computer Networks

Some Applications Using Communications Networks

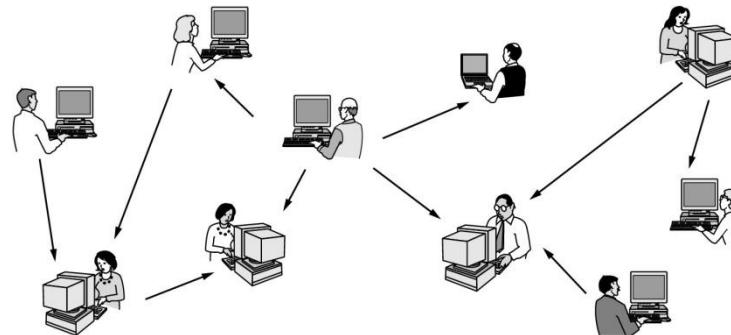
- ◆ E-mail
- ◆ Web
- ◆ Remote login
- ◆ P2P file sharing
- ◆ Multi-user network games
- ◆ Video retrieval
- ◆ Voice over IP
- ◆ Video streaming
- ◆ Real-time video conferencing
- ◆ ...

Application Architectures

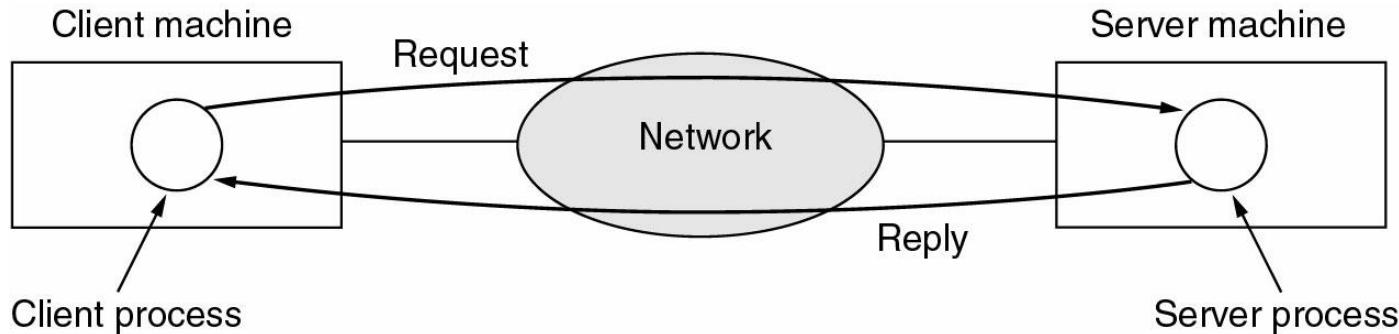
- ◆ Client-server



- ◆ Peer-to-peer (P2P)



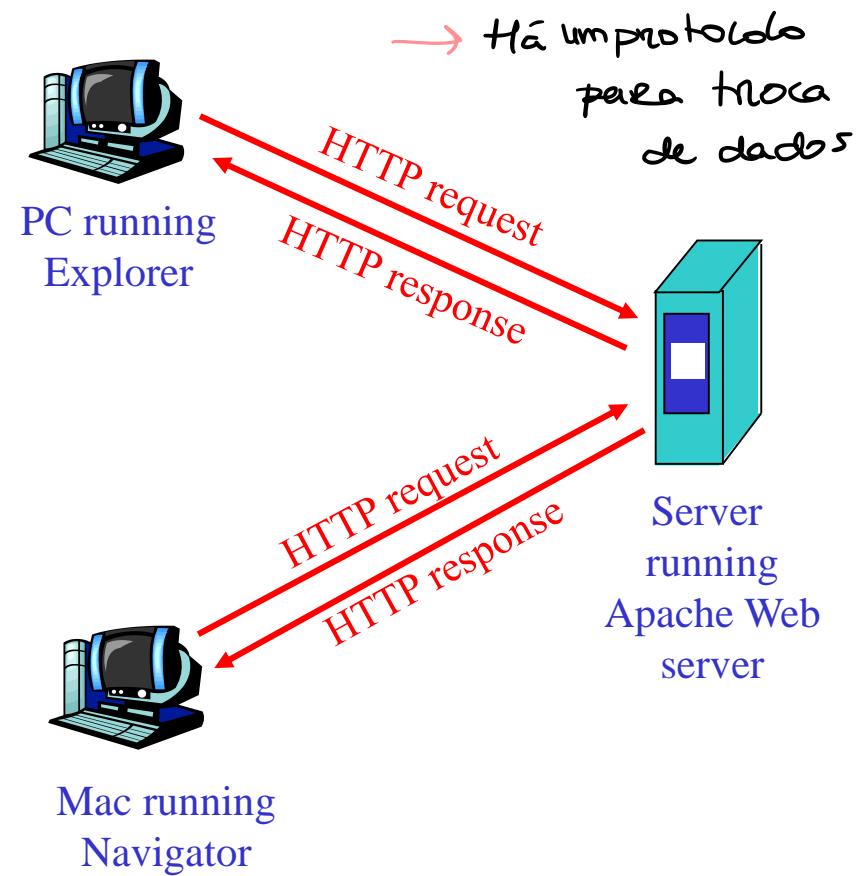
Client-server Architecture



- ◆ **Server** → Recebe os pedidos, accede à base de dados e responde
 - » always-on computer
 - » permanent IP address, well-known name
Sempre ativo
- ◆ **Clients** → Lancam processos que geram pedidos de informação ao servidor
 - » communicate with server
 - » may be intermittently connected
 - » do not communicate directly with other clients

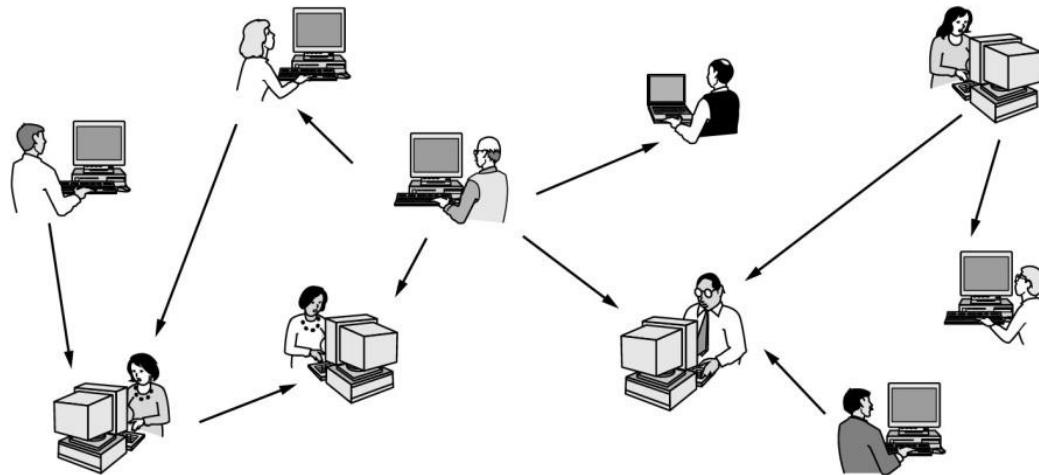
Client-server Example – The Web

- ◆ Client/server model
- ◆ Client: browser
 - » requests, receives, displays Web objects
- ◆ Server: web server
 - » sends objects in response to requests



} Nós + iguais entre si
Endereços mudam; Nós tem de estar sempre ativos

P2P Architecture



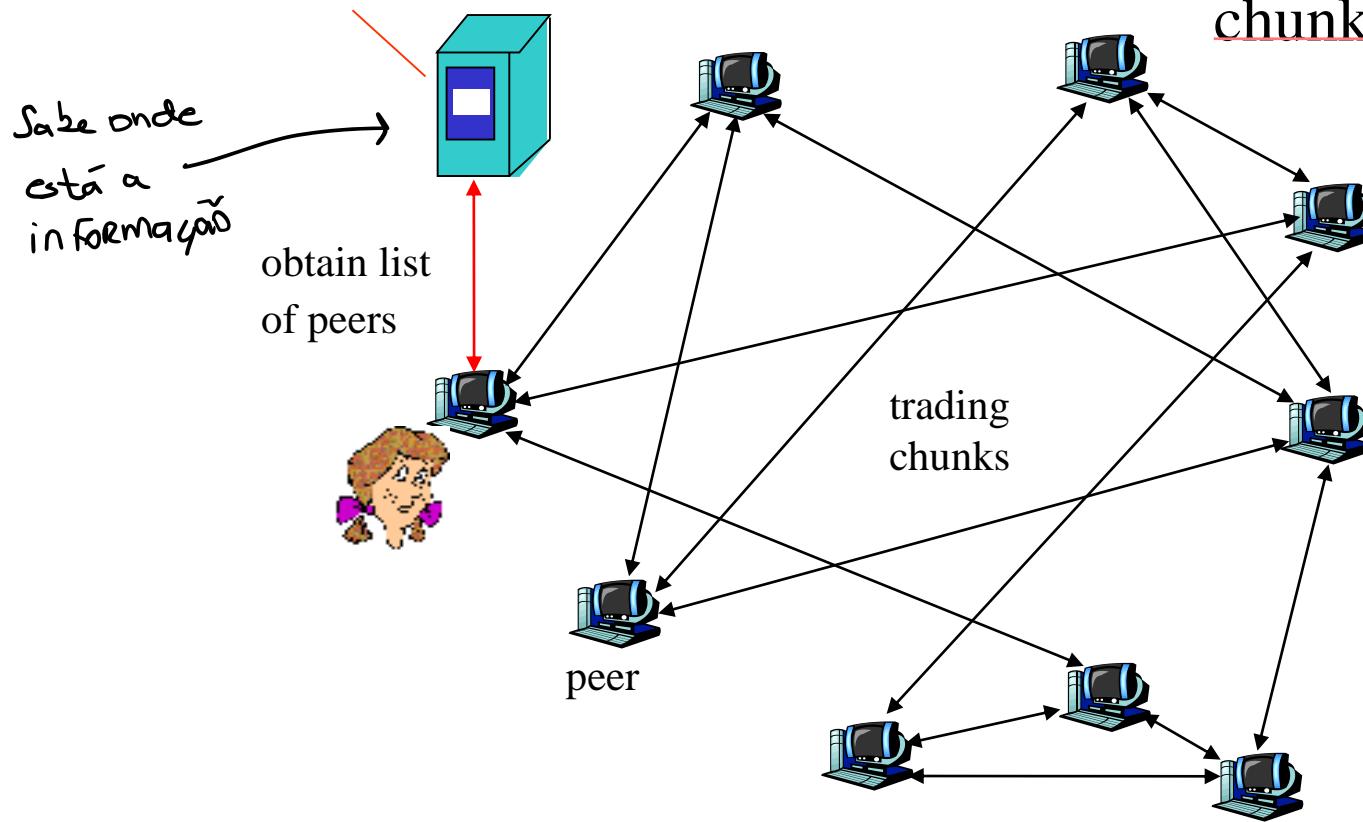
- ◆ No always-on server
- ◆ Arbitrary end systems communicate directly
- ◆ Peers are intermittently connected and may change IP addresses

P2P Example - BitTorrent

P2P file distribution

tracker: tracks peers
participating in torrent

torrent: group of
peers exchanging
chunks of a file

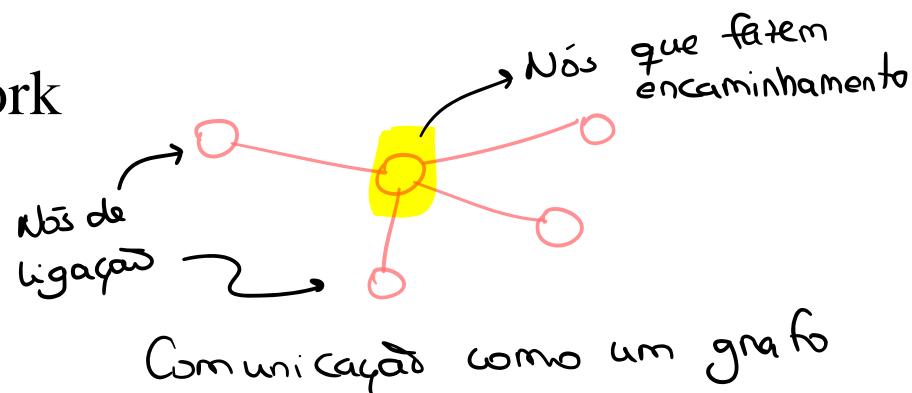


Types of Networks

Classification of Communications Networks

- ◆ By scale
 - » distance between processors
- ◆ **PAN** - Personal Area Network
- ◆ **LAN** - Local Area Network
- ◆ **MAN** - Metropolitan Area Network
- ◆ **WAN** - Wide Area Network
- ◆ **Internet**

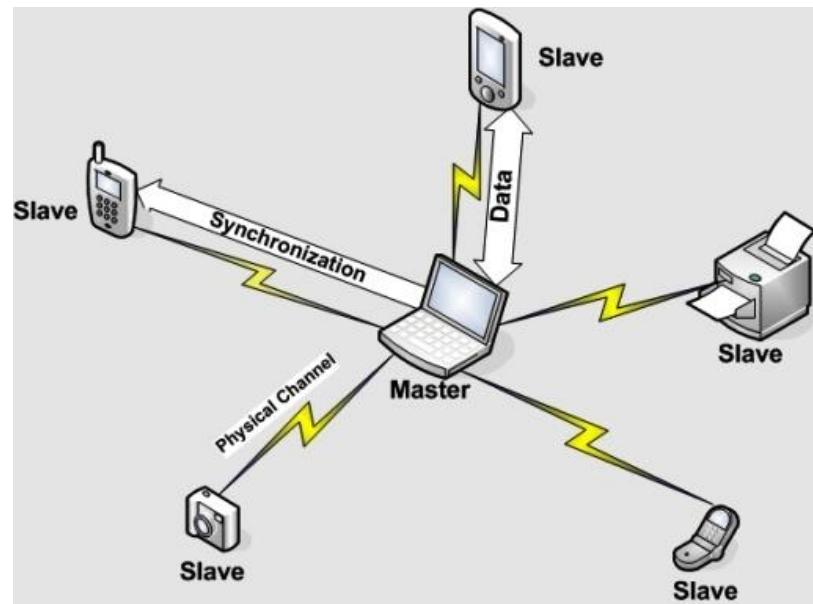
Interprocessor distance	Processors located in same	Example
1 m	Square meter	Personal area network
10 m	Room	
100 m	Building	Local area network
1 km	Campus	
10 km	City	Metropolitan area network
100 km	Country	
1000 km	Continent	Wide area network
10,000 km	Planet	The Internet



Personal Area Networks

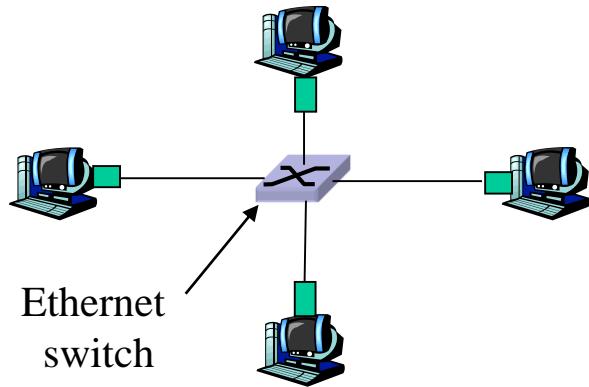
Bluetooth network

Redes pessoais à noiva
volta

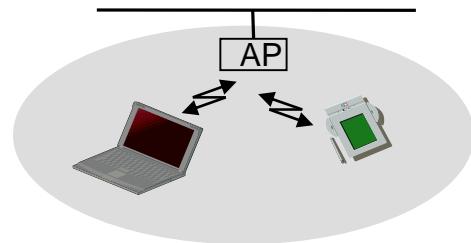


Local Area Networks

- ◆ Local Area Networks → Redes Locais

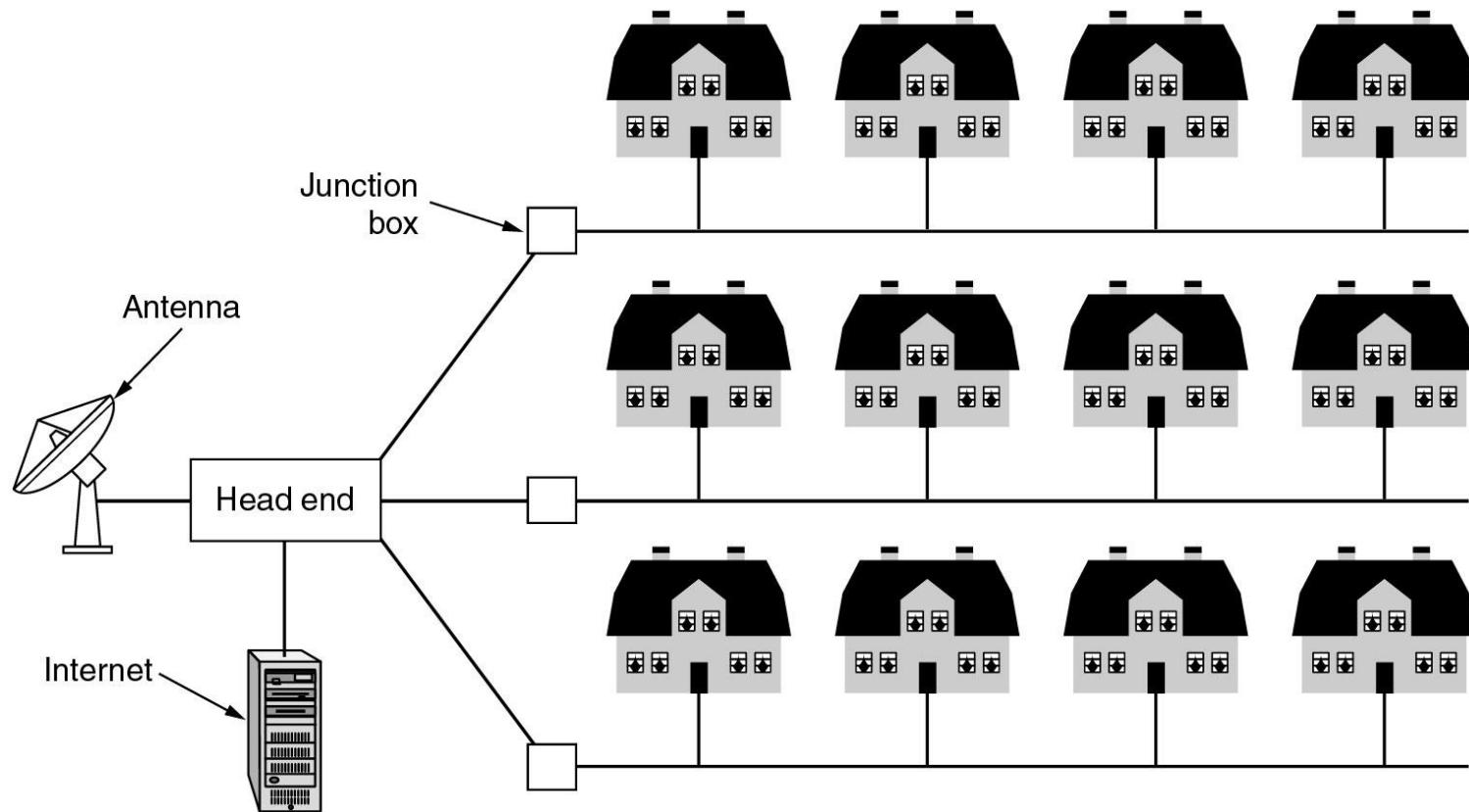


Ethernet
switch



Metropolitan Area Networks

A metropolitan area network based on cable TV



→ Cobertura

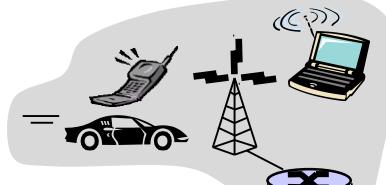
→ Rede das redes que interliga todas as redes

Internet – Interconnecting networks

◆ Network edge

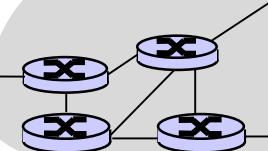
- » Hosts → Computadores telemóveis
- » Applications → Que correm nos hosts
Com necessidade de comunicação

Mobile network

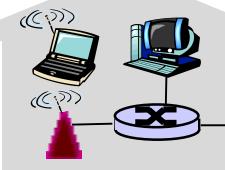


Internet Service Provider ↴

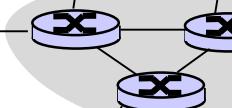
Global ISP



Home network



Regional ISP



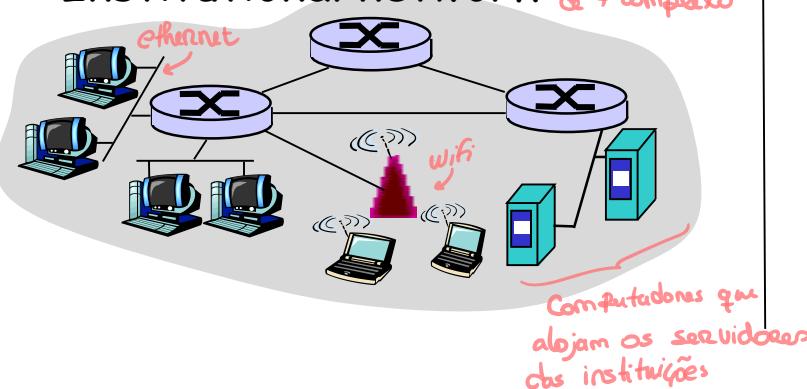
◆ Access networks

- » LANs, MANs
- » Home, Institutional
- » Mobile
- » Wired and wireless links

◆ Network core

- » Interconnected routers
- » Network of networks
- » Internet Service Providers

Institutional network



PC



server



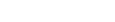
wireless laptop



cellular handheld



access points



wired links



router

ISP - Internet Service Provider

Network Software

Communications Software Organized in Black Boxes

• Dividir program
grande num +
pequeno

Passa as infoz maior
para as outras camadas
de + baixo nível

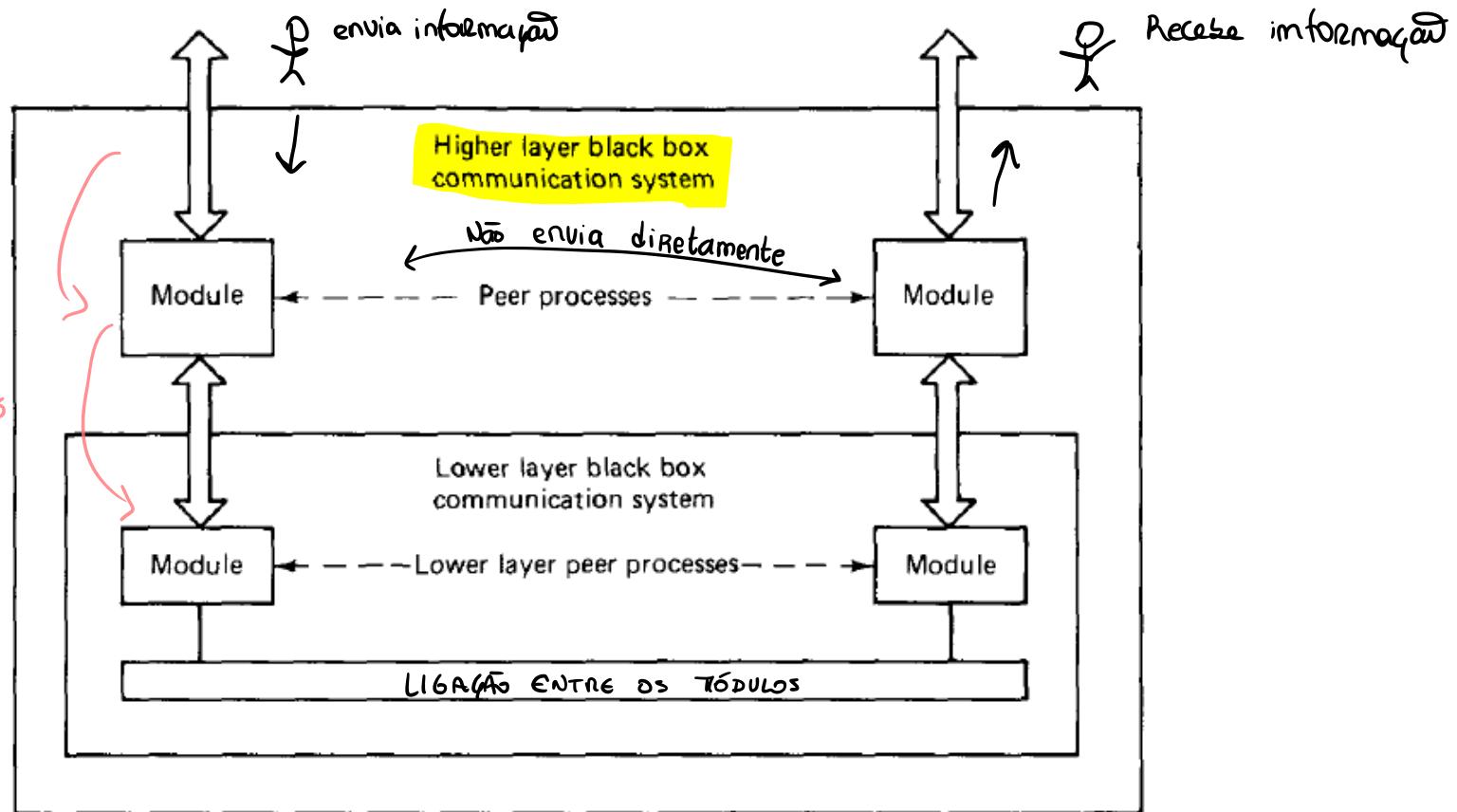
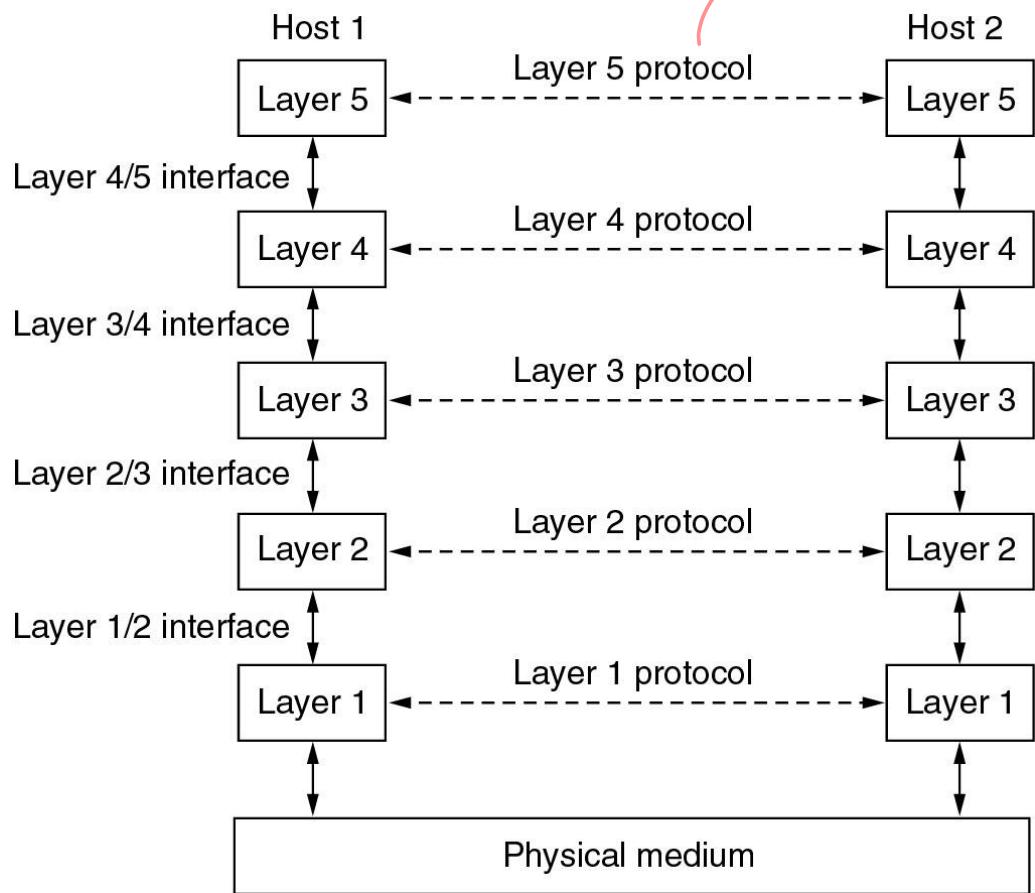


Figure 1.7 Peer processes within a black box communication system. The peer processes communicate through a lower-layer black box communication system that itself contains lower-layer peer processes.

Protocol Hierarchies

Layers, protocols, and interfaces



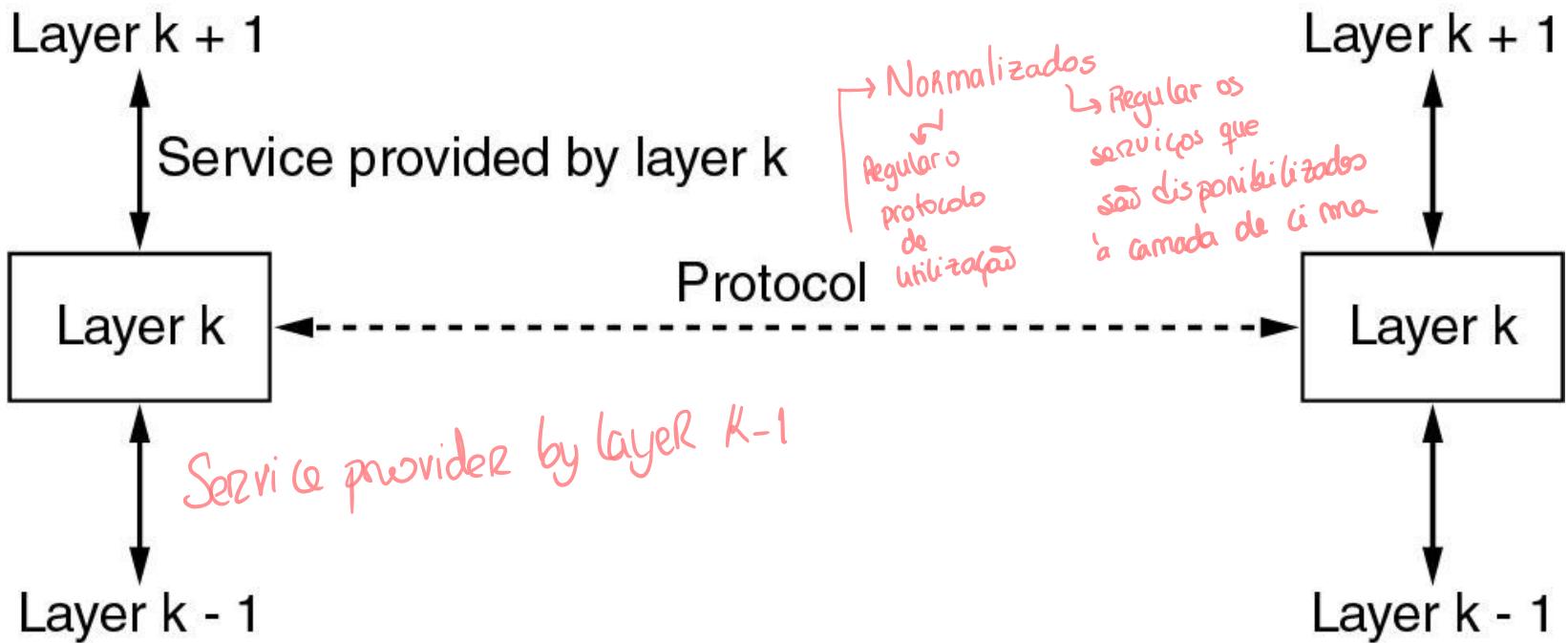
Protocolo é uma linguagem. Conjunto de mensagens para trocar a informação.

Entre módulos do mesmo nível existe um protocolo

Protocolos de diferentes níveis dizem respeito a coisas diferentes

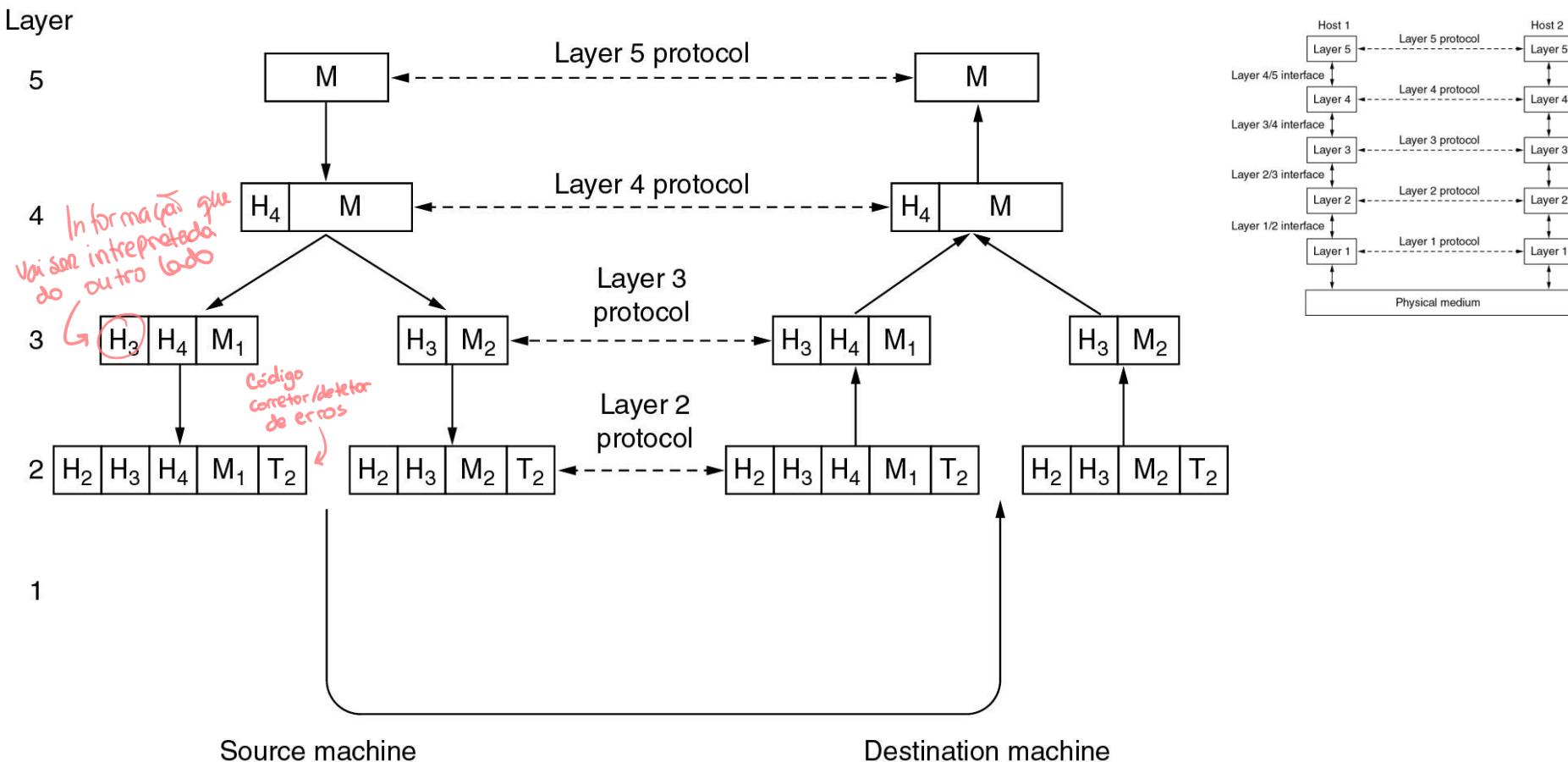
Services to Protocols Relationship

The relationship between a service and a protocol



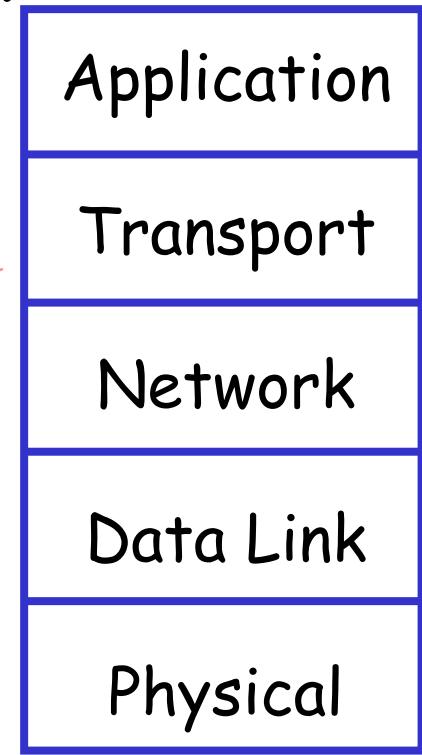
Transference of Information

Information flow supporting virtual communication in layer 5



Internet (TCP/IP) Reference Model

- ◆ Application layer ↗ Camada de + alto nível
 - » supporting network applications
 - » FTP, SMTP, HTTP, ↗ Garantir a fiabilidade na comunicação entre processos de diferentes máquinas
- ◆ Transport layer ↗ Transporte entre processos
 - » process-process (end-to-end) data transfer ↗ para garantir que não há perda de informações, congestionamento de rede...
 - » TCP, UDP
- ◆ Network layer ↗ Melhor caminho a tomar de um lado ou outro
 - » routing of data packets from source to destination
 - » IP, routing protocols
- ◆ Data Link layer → Transferir bytes de um lado para o outro com confiança
 - » data transfer between neighboring network elements
 - » PPP, Ethernet, WLAN ↗ Obter fiabilidade na comunicação entre 2 máquinas
- ◆ Physical layer
 - » bits sent “on the wire”
 - » Envio de Bits por canais/fios



Na comunicação:
↳ Usa confirmações
↳ Usa códigos } Para evitar erros

Transferring Data Through a Network

Information and Data

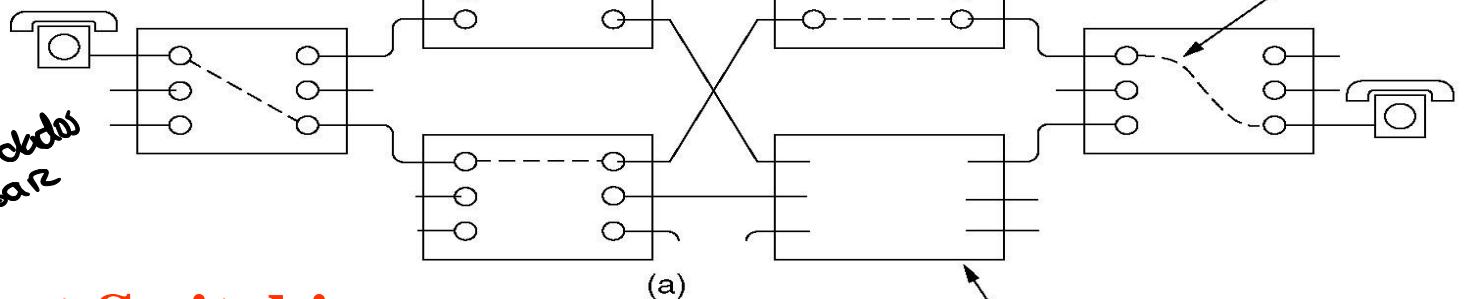
- ◆ Data
 - » term used to represent *information*
 - » e.g. text, voice, video, image, graphics
- ◆ Information represented as a sequence of bits
 - » 0110110001010....
 - » 1 Byte = 1 octet = 8 bits
 - » $1 \text{ kbit} = 10^3 \text{ bit}$; $1 \text{ Mbit} = 10^6 \text{ bit}$; $1 \text{ Gbit} = 10^9 \text{ bit}$
- ◆ Computer Networks
 - » transport information, from source to destination
 - » Information flow, capacity of a link → Byte/s; bit/s

Circuit Switching, Packet Switching

Rede que permite a comunicação entre telefones

Circuit Switching

↓
Caminho
pré-definido
por onde os dados
vão passar



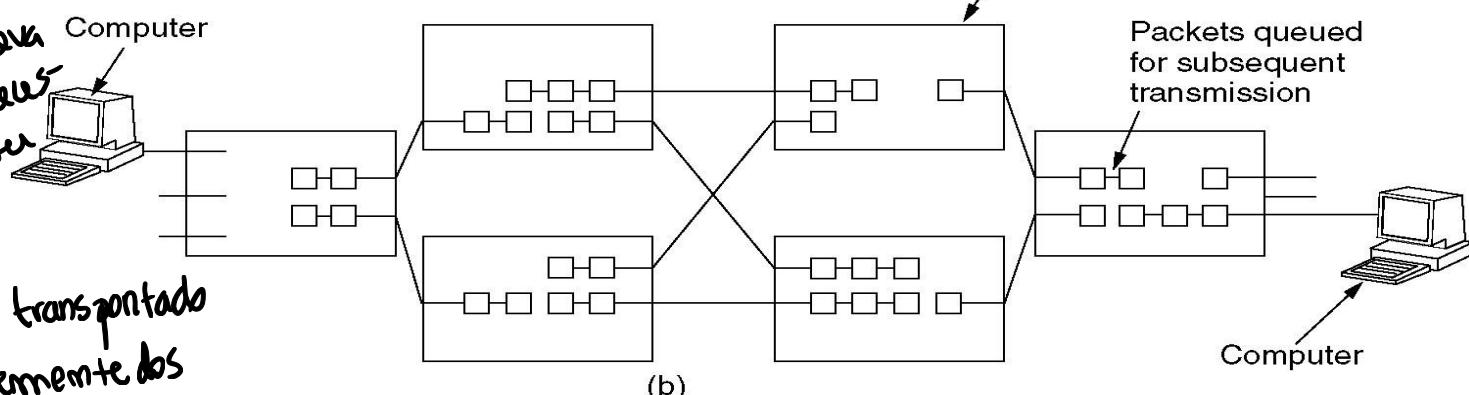
(a)

Physical copper connection set up when call is made

Packet Switching

↓
Cada pacote leva
toda a informação
sérica para ser
transferido

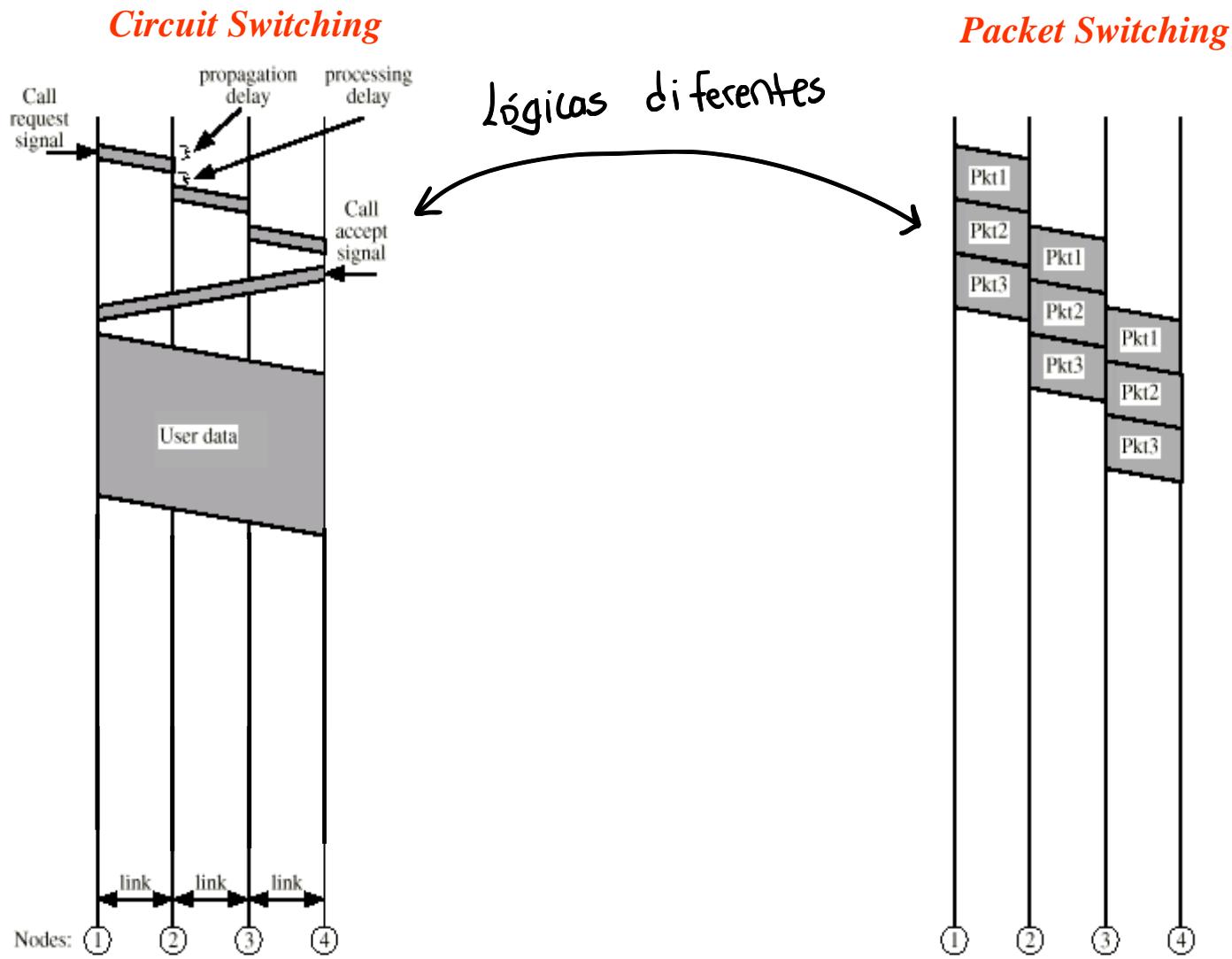
Cada pacote é transportado
independente dos
outros



(b)

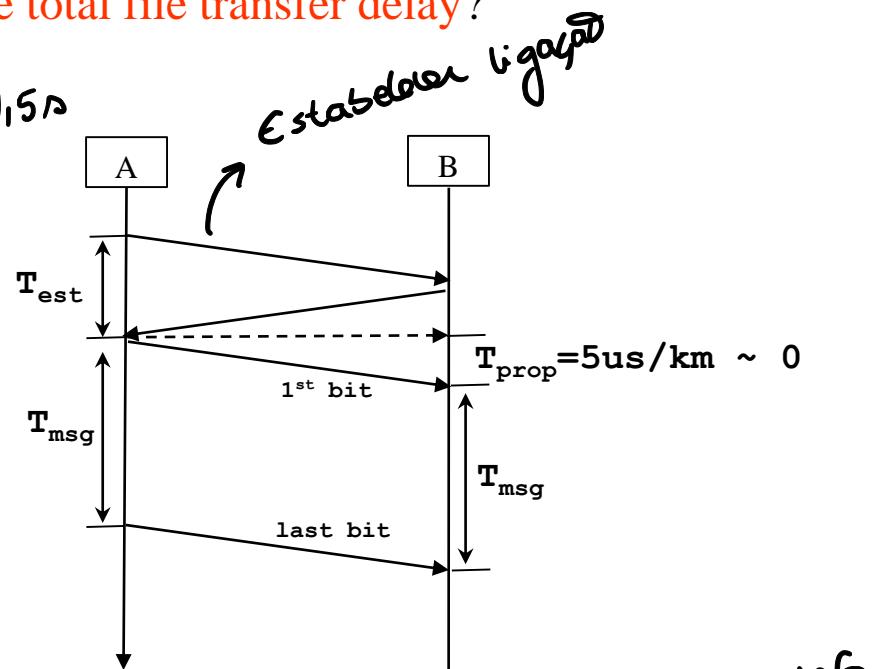
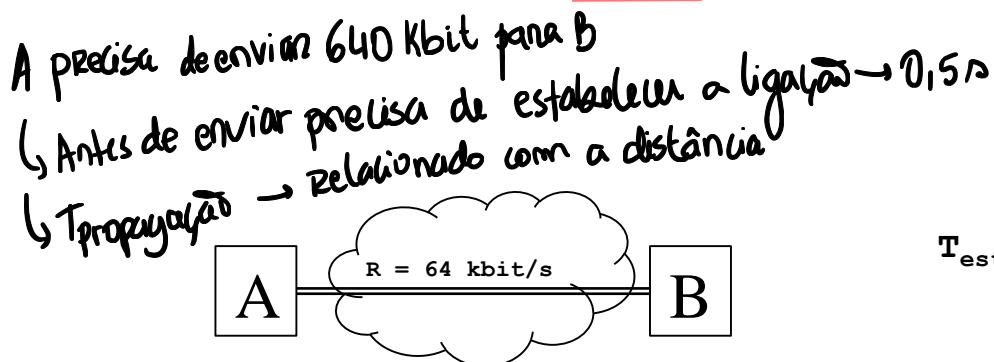
Packets queued
for subsequent
transmission

Circuit Switching, Packet Switching



Circuit Switching – Numerical Example

A file of length $L=640 \text{ kbit}$ is transferred from Host A to Host B through a circuit having a capacity of $C=64 \text{ kbit/s}$. Assuming a circuit establishment delay $T_{\text{est}}=500 \text{ ms}$, and a propagation delay $T_{\text{prop}} \sim 0$, what is the total file transfer delay?



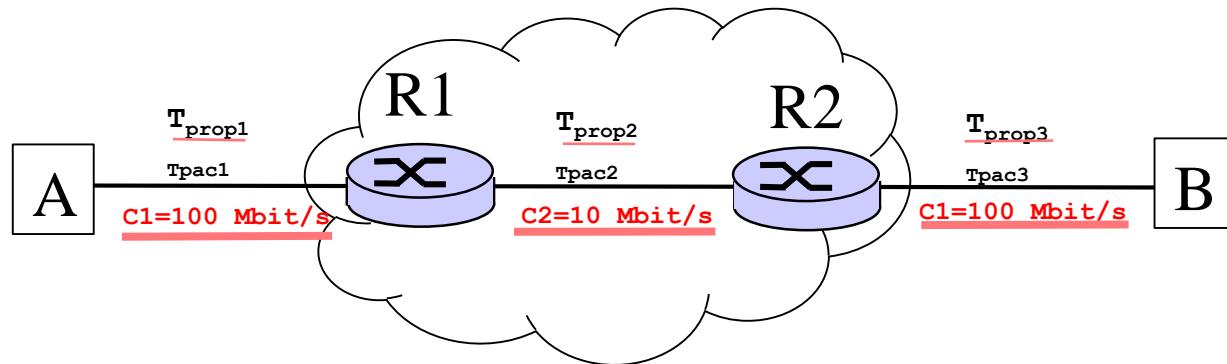
◆ Answer:

» $T_{\text{msg}} = L/C = 640 \text{ kbit} / 64 \text{ kbit/s} = 10 \text{ s} \rightarrow 10 \text{ s para transmitir info}$

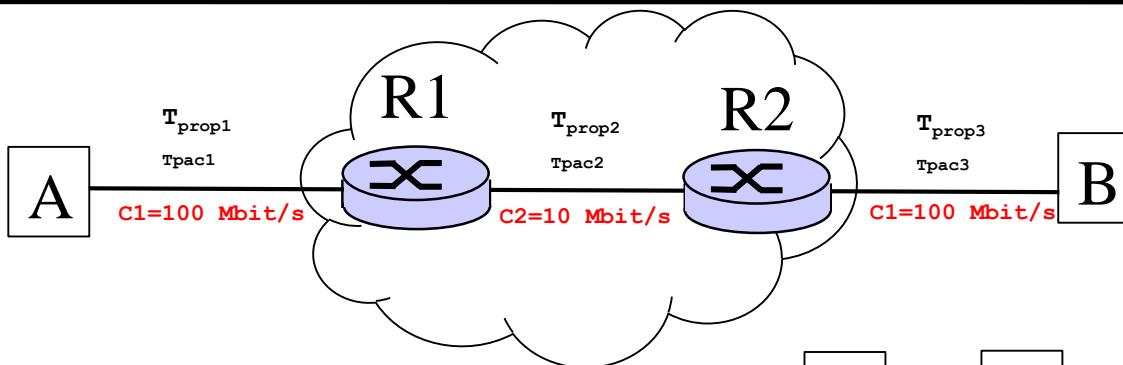
» $T_{\text{tot}} = T_{\text{est}} + T_{\text{prop}} + T_{\text{msg}} = 0.5 + 0 + 10 = 10.5 \text{ s}$

Packet Switching – Numerical Example

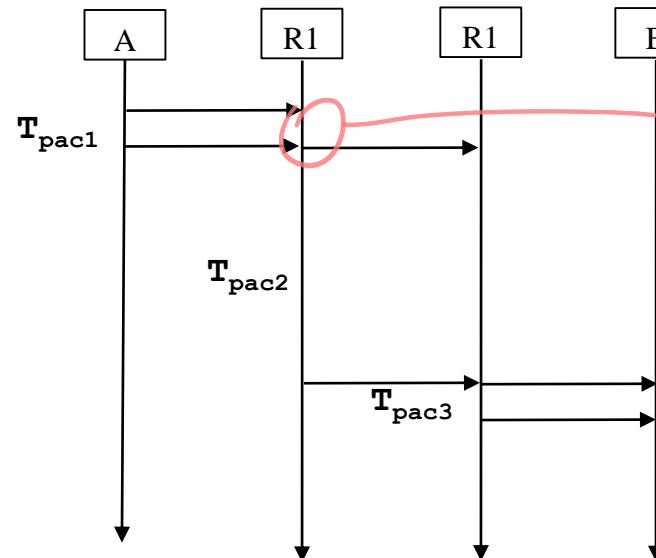
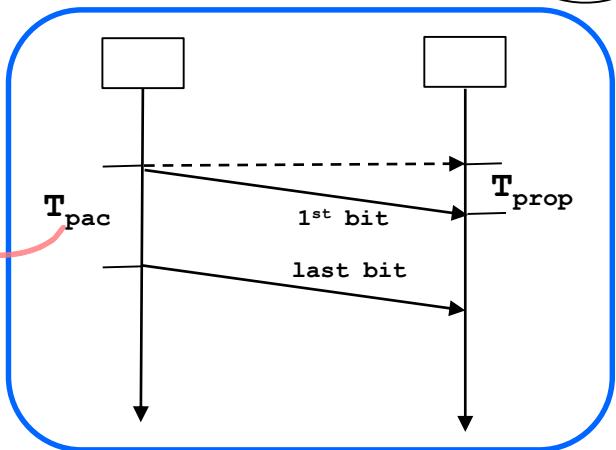
Host A sends a packet of length L=10 kbit to Host B through routers R1 and R2. Assuming propagation delay through the 3 links is T_{prop}~0 and that there are no queuing delays at the network elements (A, R1 and R2), what is the end-to-end packet delay?



Packet Switching – Numerical Example



tempo desde que mete o primeiro bit até que chega ao destinatário
 $T_{prop} = 0$
 $T_{queue} = 0$



Answer:

» $T_{pac1} = T_{pac3} = L/C_1 = 10 \text{ kbit} / 100 \text{ Mbit/s} = 0.1 \text{ ms}$

» $T_{pac2} = L/C_2 = 10 \text{ kbit} / 10 \text{ Mbit/s} = 1 \text{ ms}$

» $T_{end-to-end} = T_{pac1} + T_{prop1} + T_{pac2} + T_{prop2} + T_{pac3} + T_{prop3} = 1.2 \text{ ms}$

Homework

1. Review slides
2. Read from Tanenbaum
 - » Chapter 1 - Introduction
 - » Section 2.6.5 – Switching
3. Answer questions at moodle