

# Network Security - Week 4

Manuel E. Correia

DCC/FCUP

2025

# Alternative authentication - Kerberos

## In Greek mythology

- A ferocious 3-headed dog
- Guards the entrance to Hades' realm



# Alternative authentication - Kerberos

## In Greek mythology

- A ferocious 3-headed dog
- Guards the entrance to Hades' realm



## In Security

- An authentication protocol
- Guards the entrance to some realms
- Designed for smaller scale, e.g. LANs
- Relies on a Trusted Third Party (TTP)



# Motivation for Kerberos

- Authentication using public keys
  - $N$  users  $\Rightarrow N$  key pairs
  - Slow

# Motivation for Kerberos

- Authentication using public keys
  - $N$  users  $\Rightarrow N$  key pairs
  - Slow
- Authentication using symmetric keys
  - Fast
  - $N$  users requires  $N^2$  keys (roughly)

# Motivation for Kerberos

- Authentication using public keys
  - $N$  users  $\Rightarrow N$  key pairs
  - Slow
- Authentication using symmetric keys
  - Fast
  - $N$  users requires  $N^2$  keys (roughly)
- Symmetric key case **does not scale**

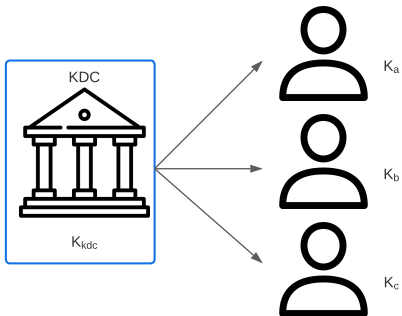
# Motivation for Kerberos

- Authentication using public keys
  - $N$  users  $\Rightarrow$   $N$  key pairs
  - Slow
- Authentication using symmetric keys
  - Fast
  - $N$  users requires  $N^2$  keys (roughly)
- Symmetric key case **does not scale**

Kerberos is based on symmetric keys, but only requires  $N$  keys.

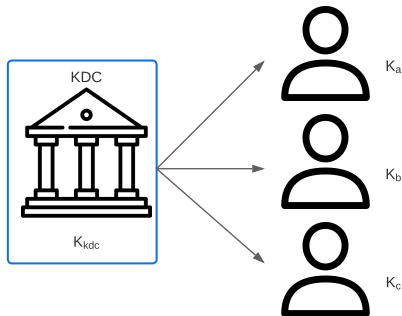
- *Assumption:* Security depends on TTP
- No PKI necessary!

# A Trusted Key Distribution Center



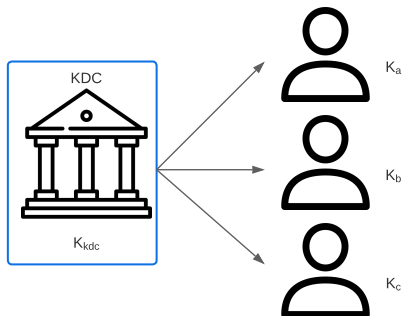


# A Trusted Key Distribution Center



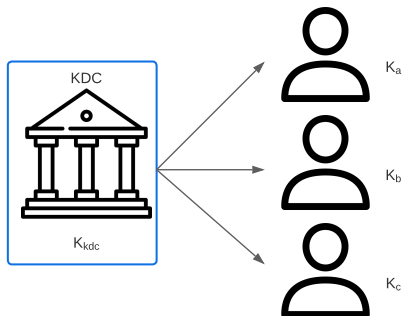
- KDC acts as a TTP (CA in PKIs)
- Cannot be compromised (by assumption)

# A Trusted Key Distribution Center



- KDC acts as a TTP (CA in PKIs)
- Cannot be compromised (by assumption)
- Shares symmetric keys with all participants ( $K_a, K_b, K_c$ )
- ... and an extra master key  $K_{kdc}$

# A Trusted Key Distribution Center



- KDC acts as a TTP (CA in PKIs)
- Cannot be compromised (by assumption)
- Shares symmetric keys with all participants ( $K_a, K_b, K_c$ )
- ... and an extra master key  $K_{kdc}$
- Enables authentication and session keys

## Tickets within tickets

- KDC grants tickets needed to access network resources

## Tickets within tickets

- KDC grants tickets needed to access network resources
- But also grants Ticket-Granting Tickets (TGTs)
  - Issued upon first login
  - Used to obtain regular tickets

# Kerberos Authentication

## Tickets within tickets

- KDC grants tickets needed to access network resources
- But also grants Ticket-Granting Tickets (TGTs)
  - Issued upon first login
  - Used to obtain regular tickets

## TGTs

- Each TGT contains:
  - Session key
  - User ID
  - Expiration time

# Kerberos Authentication

## Tickets within tickets

- KDC grants tickets needed to access network resources
- But also grants Ticket-Granting Tickets (TGTs)
  - Issued upon first login
  - Used to obtain regular tickets

## TGTs

- Each TGT contains:
  - Session key
  - User ID
  - Expiration time
- Used to avoid having the KDC manage a database
- KDC remains (mostly) stateless

# Kerberos Login

- 1 Alice enters her password



# Kerberos Login

- 1 Alice enters her password
- 2 Alice retrieves the TGT
  - Derives  $K_A$  from its password.
  - Uses  $K_A$  to request TGT from the KDC

# Kerberos Login

- 1 Alice enters her password
- 2 Alice retrieves the TGT
  - Derives  $K_A$  from its password.
  - Uses  $K_A$  to request TGT from the KDC
- 3 Forward the TGT to access network resources

# Kerberos Login

- 1 Alice enters her password
- 2 Alice retrieves the TGT
  - Derives  $K_A$  from its password.
  - Uses  $K_A$  to request TGT from the KDC
- 3 Forward the TGT to access network resources

The good and the bad

# Kerberos Login

- 1 Alice enters her password
- 2 Alice retrieves the TGT
  - Derives  $K_A$  from its password.
  - Uses  $K_A$  to request TGT from the KDC
- 3 Forward the TGT to access network resources

## The good and the bad

- **Good:** Security is transparent for Alice

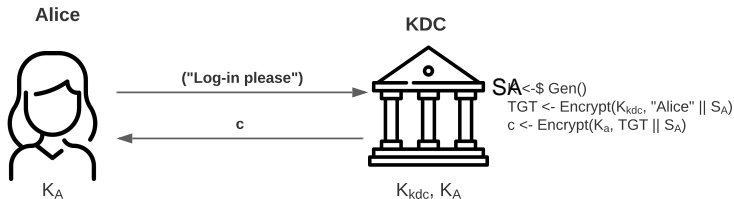
# Kerberos Login

- 1 Alice enters her password
- 2 Alice retrieves the TGT
  - Derives  $K_A$  from its password.
  - Uses  $K_A$  to request TGT from the KDC
- 3 Forward the TGT to access network resources

## The good and the bad

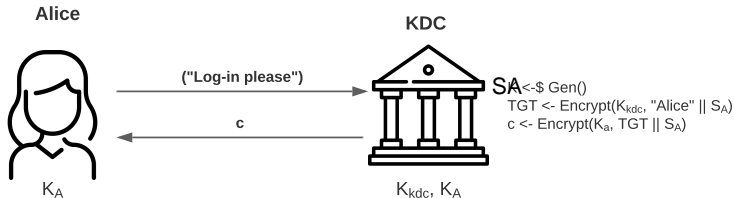
- **Good:** Security is transparent for Alice
- **Bad:** KDC must always be secure!

# Kerberos Login



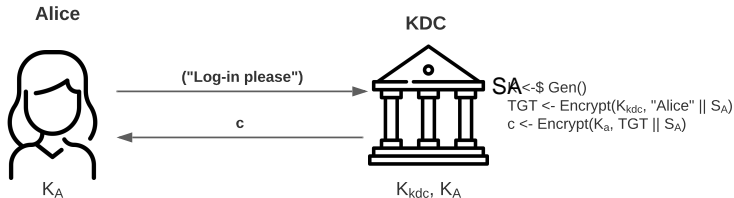
- $K_A$  is Alice's key (e.g. generated from password)

# Kerberos Login



- $K_A$  is Alice's key (e.g. generated from password)
- KDC generates session key  $S_A$  and constructs TGT

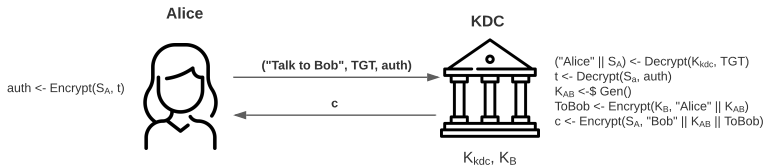
# Kerberos Login



- $K_A$  is Alice's key (e.g. generated from password)
- KDC generates session key  $S_A$  and constructs TGT
- Alice retrieves  $S_A$  and TGT

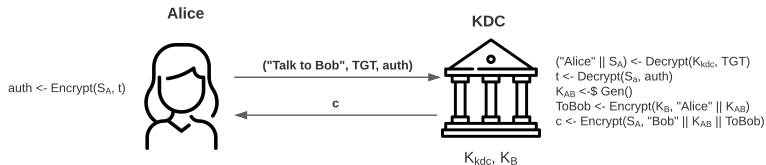


# Kerberos Talk to Bob - P1



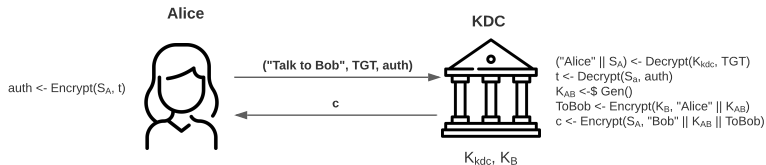
- KDC knows Bob's key  $K_B$

# Kerberos Talk to Bob - P1



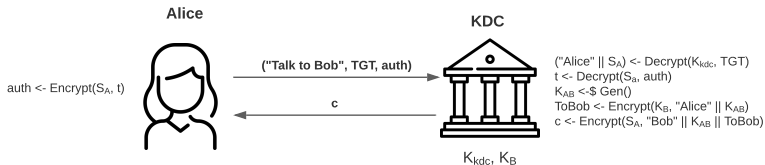
- KDC knows Bob's key  $K_B$
- Alice sends the request, alongside the TGT and an authenticator

# Kerberos Talk to Bob - P1



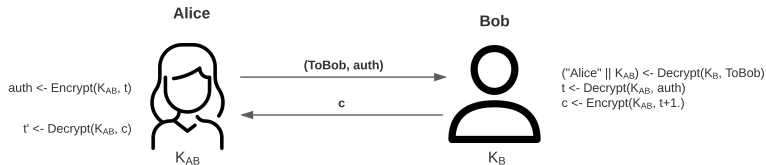
- KDC knows Bob's key  $K_B$
- Alice sends the request, alongside the TGT and an authenticator
- KDC prepares a communication key for  $K_{AB}$ 
  - Encrypts it also with  $K_B$
  - And tags Alice (to avoid reflection attacks)

# Kerberos Talk to Bob - P1



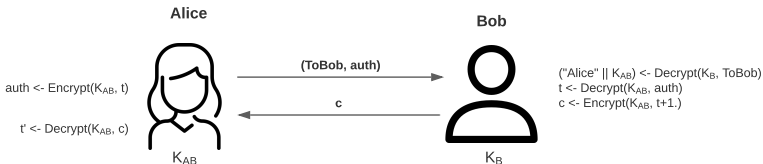
- KDC knows Bob's key  $K_B$
- Alice sends the request, alongside the TGT and an authenticator
- KDC prepares a communication key for  $K_{AB}$ 
  - Encrypts it also with  $K_B$
  - And tags Alice (to avoid reflection attacks)
- Alice retrieves  $K_{AB}$  and an authenticator it can send to Bob

# Kerberos Talk to Bob - P2



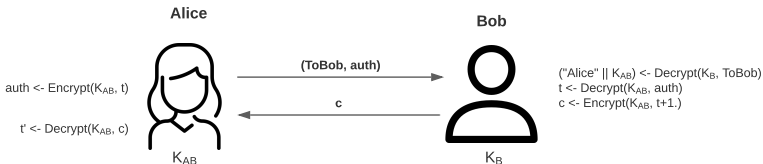
- Bob knows its own key  $K_B$

# Kerberos Talk to Bob - P2



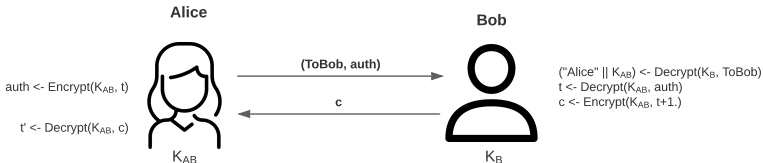
- Bob knows its own key  $K_B$
- Alice sends ToBob, and an authenticator encrypted with  $K_{AB}$

# Kerberos Talk to Bob - P2



- Bob knows its own key  $K_B$
- Alice sends ToBob, and an authenticator encrypted with  $K_{AB}$
- Bob does not know  $K_{AB}$ ...
  - But the ToBob token has  $K_{AB}$ , encrypted with  $K_B$
  - Retrieves  $K_{AB}$  and checks the authenticator for freshness
  - Encrypts a reply with the updated timestamp

# Kerberos Talk to Bob - P2



- Bob knows its own key  $K_B$
- Alice sends ToBob, and an authenticator encrypted with  $K_{AB}$
- Bob does not know  $K_{AB}$ ...
  - But the ToBob token has  $K_{AB}$ , encrypted with  $K_B$
  - Retrieves  $K_{AB}$  and checks the authenticator for freshness
  - Encrypts a reply with the updated timestamp
- Alice decrypts the reply and checks for freshness



- Key  $S_A$  is used for authentication
  - Gives confidentiality/integrity for Alice-KDC communication

- Key  $S_A$  is used for authentication
  - Gives confidentiality/integrity for Alice-KDC communication
- Key  $K_{AB}$  used for Alice-Bob communication
  - Trustworthiness from the fact that the KDC encrypt it with  $K_B$
  - Only entity with knowledge of  $K_B$
  - Bob trusts the KDC!

- Key  $S_A$  is used for authentication
  - Gives confidentiality/integrity for Alice-KDC communication
- Key  $K_{AB}$  used for Alice-Bob communication
  - Trustworthiness from the fact that the KDC encrypt it with  $K_B$
  - Only entity with knowledge of  $K_B$
  - **Bob trusts the KDC!**
- Timestamps are used for authentication and replay protection

- Key  $S_A$  is used for authentication
  - Gives confidentiality/integrity for Alice-KDC communication
- Key  $K_{AB}$  used for Alice-Bob communication
  - Trustworthiness from the fact that the KDC encrypt it with  $K_B$
  - Only entity with knowledge of  $K_B$
  - **Bob trusts the KDC!**
- Timestamps are used for authentication and replay protection
- Timestamps behave like a nonce that is known in advance

- Key  $S_A$  is used for authentication
  - Gives confidentiality/integrity for Alice-KDC communication
- Key  $K_{AB}$  used for Alice-Bob communication
  - Trustworthiness from the fact that the KDC encrypt it with  $K_B$
  - Only entity with knowledge of  $K_B$
  - **Bob trusts the KDC!**
- Timestamps are used for authentication and replay protection
- Timestamps behave like a nonce that is known in advance
- “time” is a security-critical parameter!

## Common questions

- Why is TGT encrypted with  $K_A$ ?
  - Allows for the anonymity of Alice!

## Common questions

- Why is TGT encrypted with  $K_A$ ?
  - Allows for the anonymity of Alice!
- Why doesn't the KDC sent the ticket to Bob?
  - Bob would have to remember  $K_{AB}$
  - It's good to remain stateless!

## Common questions

- Why is TGT encrypted with  $K_A$ ?
  - Allows for the anonymity of Alice!
- Why doesn't the KDC send the ticket to Bob?
  - Bob would have to remember  $K_{AB}$
  - It's good to remain stateless!
- Can't we have the KDC remember the session key instead of using the TGT?
  - Yes... but it's good to have the KDC remain stateless
  - Scales better!



# Wrap up

## Authentication is core in network security

- Single or mutual authentication

# Wrap up

## Authentication is core in network security

- Single or mutual authentication
- Vulnerable to replay, relay, mitm attacks

# Wrap up

## Authentication is core in network security

- Single or mutual authentication
- Vulnerable to replay, relay, mitm attacks
- Nonces/timestamps used for freshness

# Wrap up

## Authentication is core in network security

- Single or mutual authentication
- Vulnerable to replay, relay, mitm attacks
- Nonces/timestamps used for freshness

## Cryptogaphy

- Signatures used to ensure source is trustworthy

# Wrap up

## Authentication is core in network security

- Single or mutual authentication
- Vulnerable to replay, relay, mitm attacks
- Nonces/timestamps used for freshness

## Cryptogaphy

- Signatures used to ensure source is trustworthy
- Encryption used to ensure confidentiality

# Wrap up

## Authentication is core in network security

- Single or mutual authentication
- Vulnerable to replay, relay, mitm attacks
- Nonces/timestamps used for freshness

## Cryptogaphy

- Signatures used to ensure source is trustworthy
- Encryption used to ensure confidentiality

## Kerberos

- An alternative system for authentication

# Wrap up

## Authentication is core in network security

- Single or mutual authentication
- Vulnerable to replay, relay, mitm attacks
- Nonces/timestamps used for freshness

## Cryptogaphy

- Signatures used to ensure source is trustworthy
- Encryption used to ensure confidentiality

## Kerberos

- An alternative system for authentication
- Trusted hardware assumption - KDC

# Wrap up

## Authentication is core in network security

- Single or mutual authentication
- Vulnerable to replay, relay, mitm attacks
- Nonces/timestamps used for freshness

## Cryptogaphy

- Signatures used to ensure source is trustworthy
- Encryption used to ensure confidentiality

## Kerberos

- An alternative system for authentication
- Trusted hardware assumption - KDC
- Ticket Granting Tickets
  - Allows for stateless resource management



# Network Security - Week 4

Manuel E. Correia

DCC/FCUP

2025