

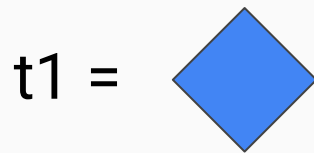
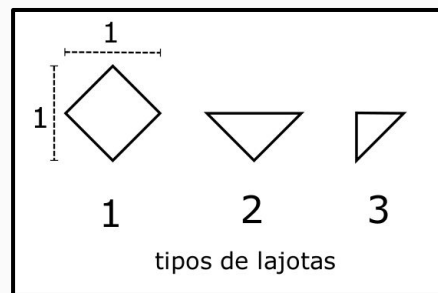
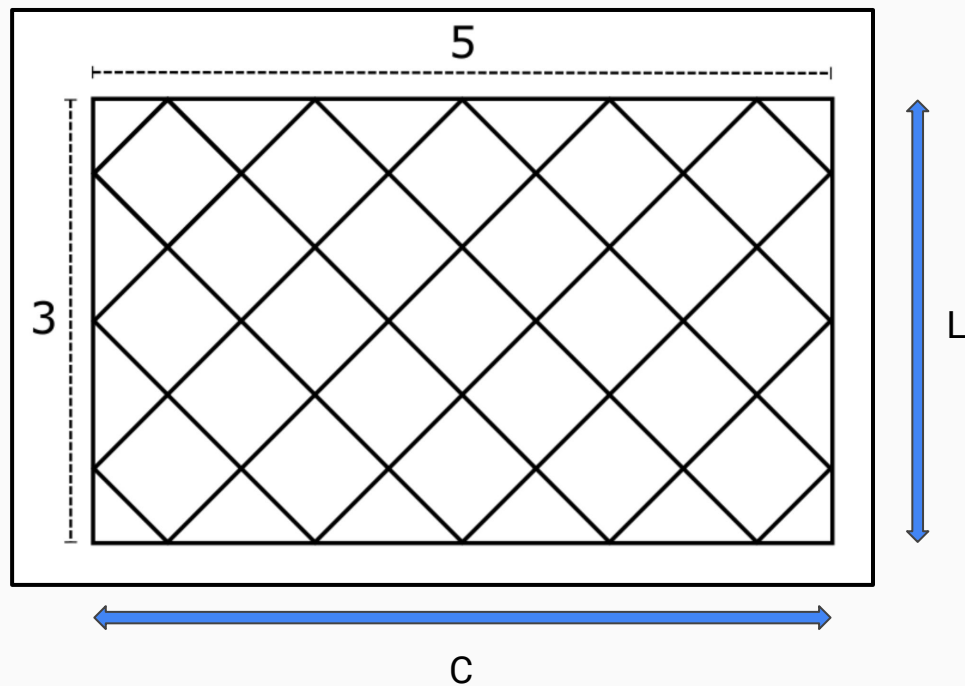


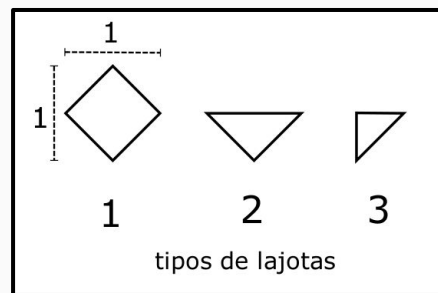
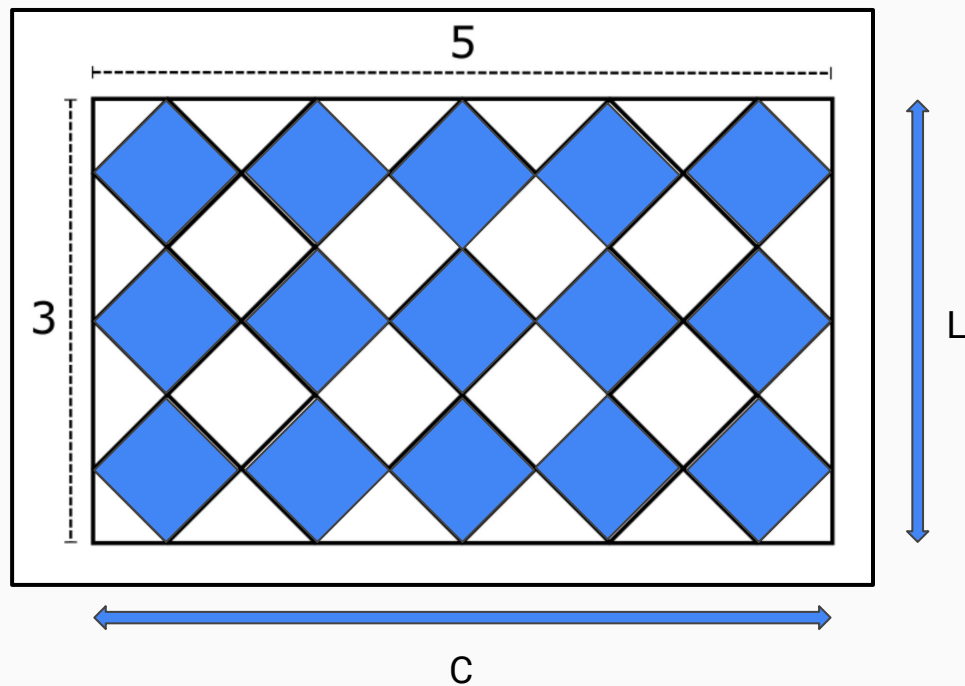
OBI2018

XX Olimpíada Brasileira de Informática

Piso da escola

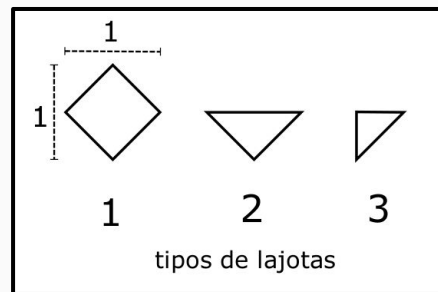
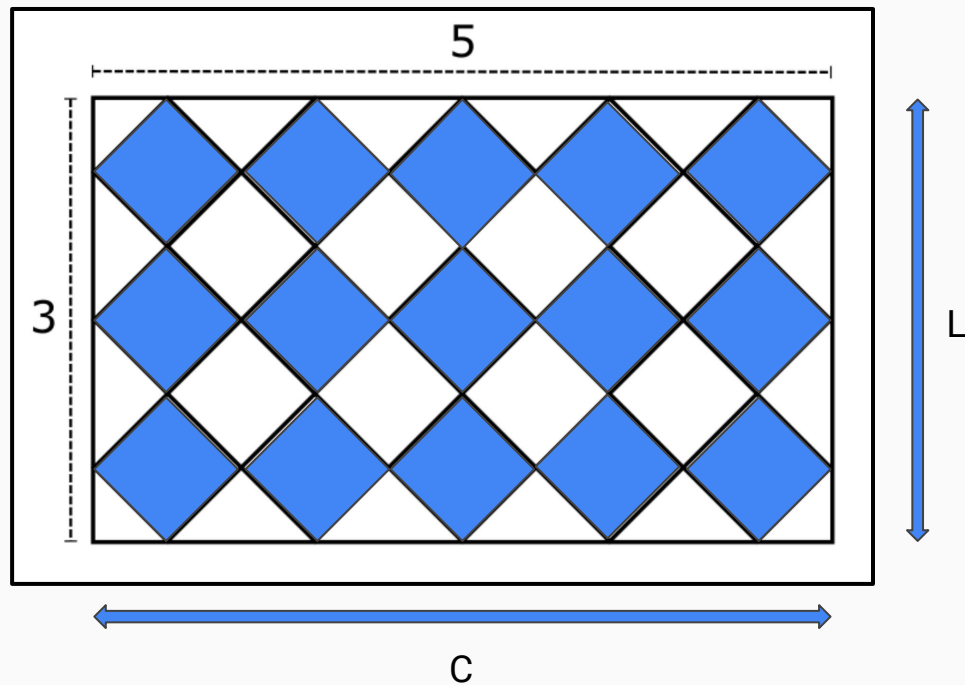
- Fácil
- Entrada e saída
- Matemática





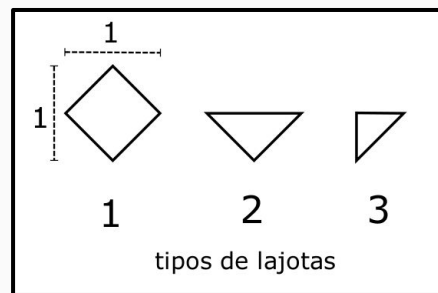
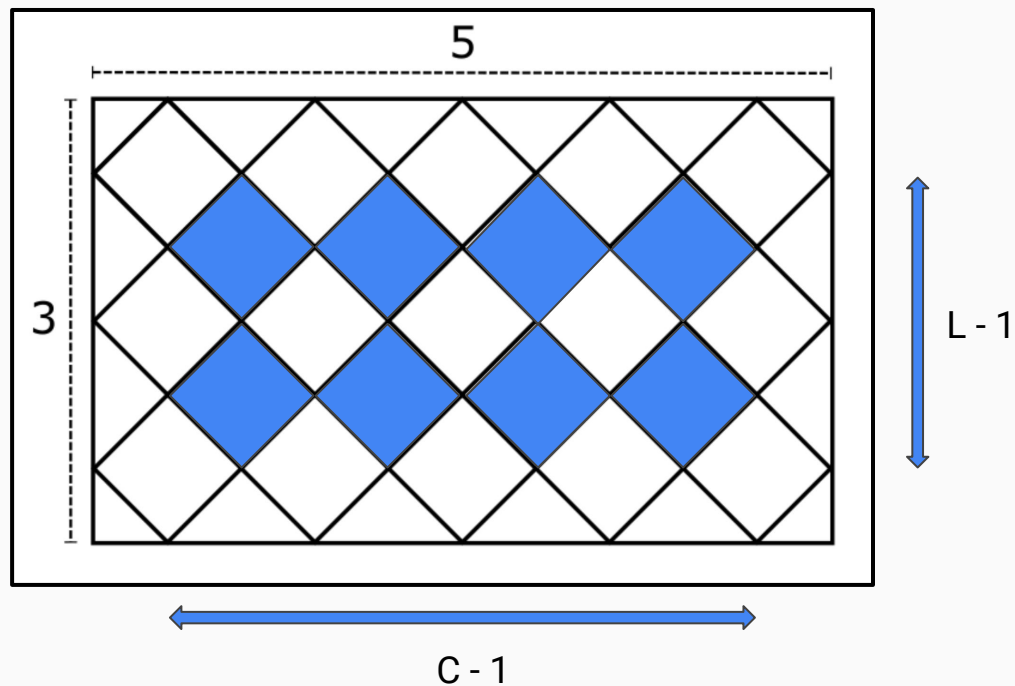
$t_1 =$

$t_2 =$



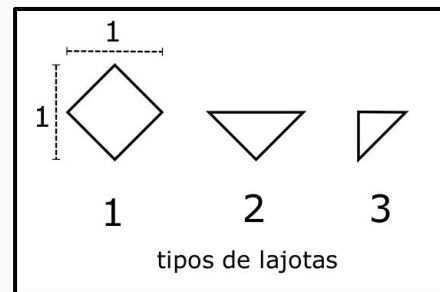
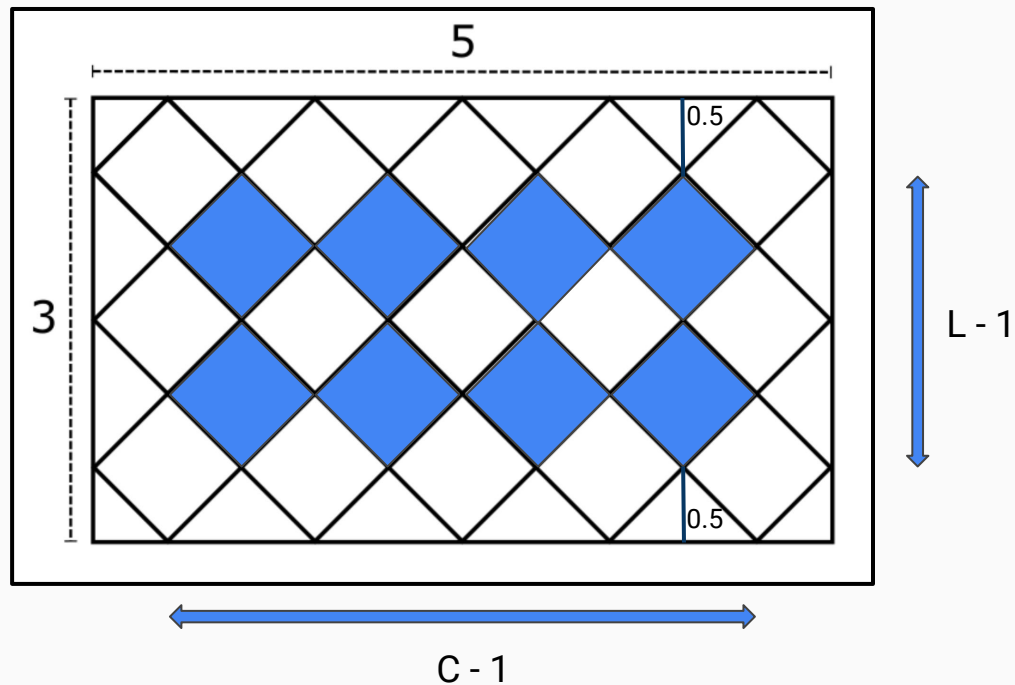
$$t1 = L * C + ?$$

$$t2 =$$



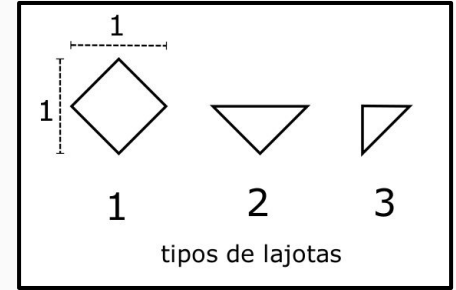
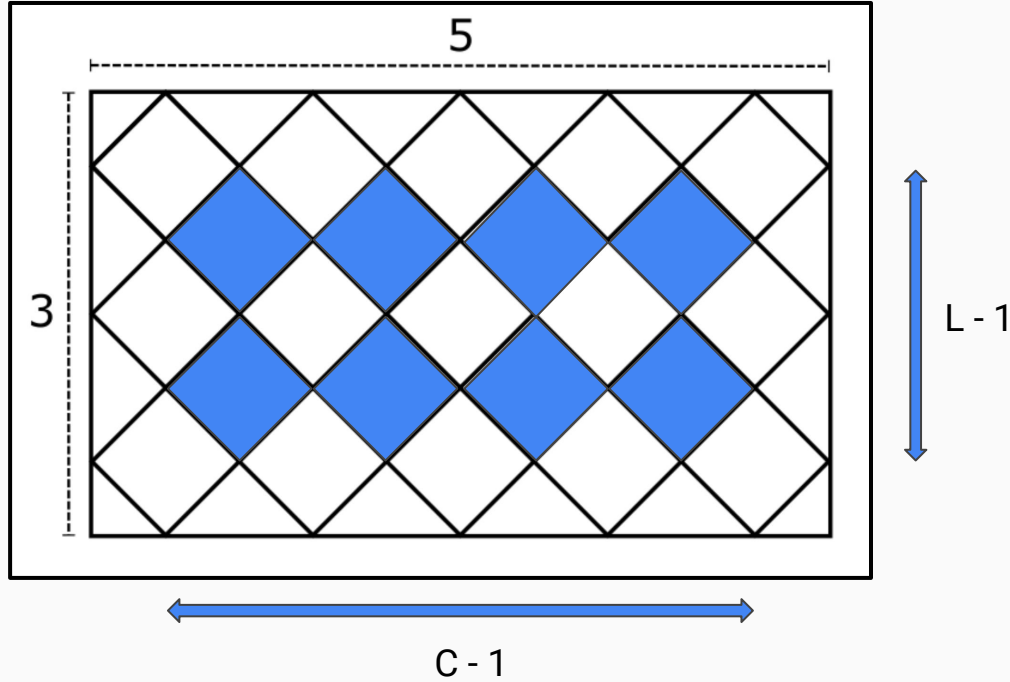
$$t1 = L * C + ?$$

$$t2 =$$



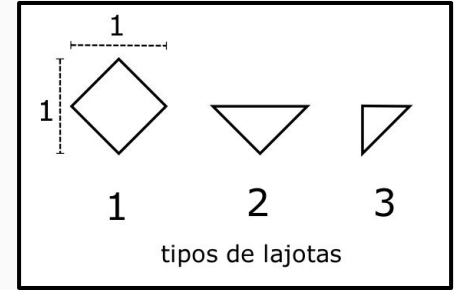
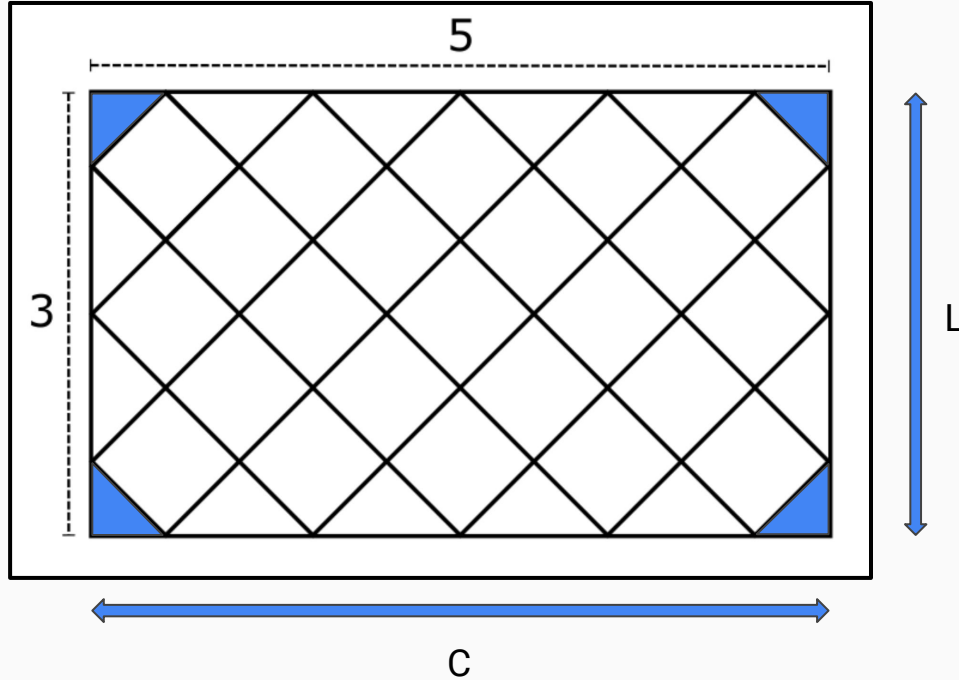
$$t1 = L * C + ?$$

$$t2 =$$



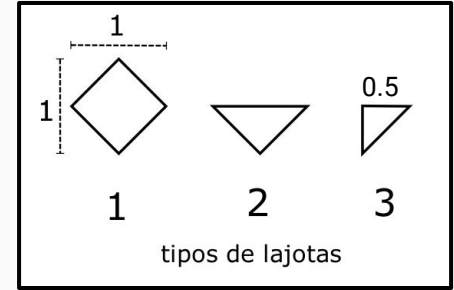
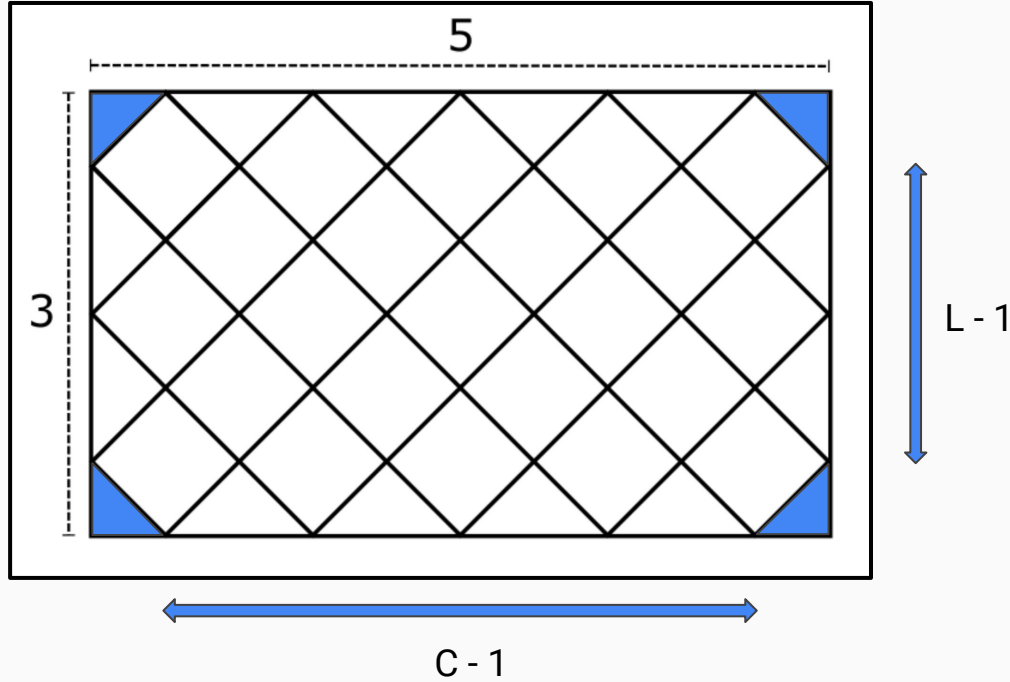
$$t1 = L * C + (L - 1)(C - 1)$$

$$t2 =$$



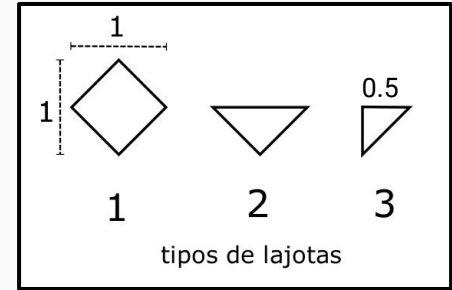
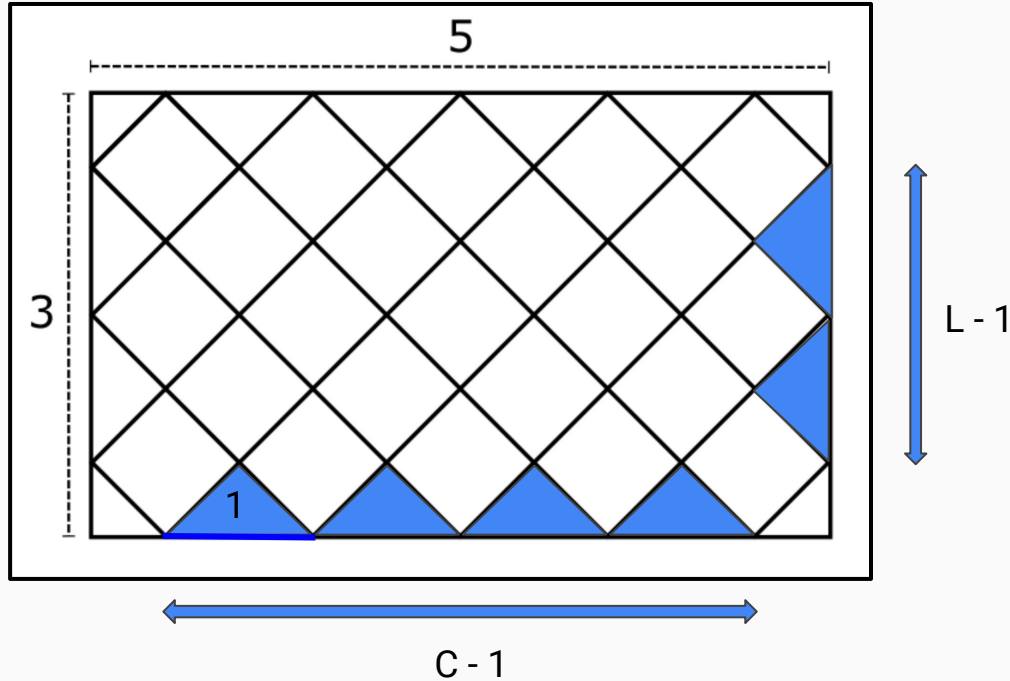
$$t1 = L * C + (L - 1)(C - 1)$$

$$t2 =$$



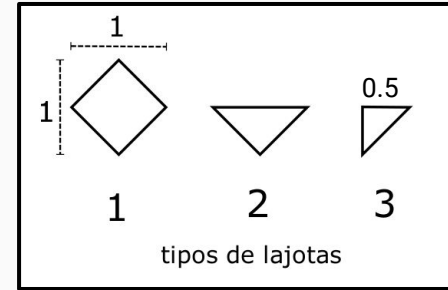
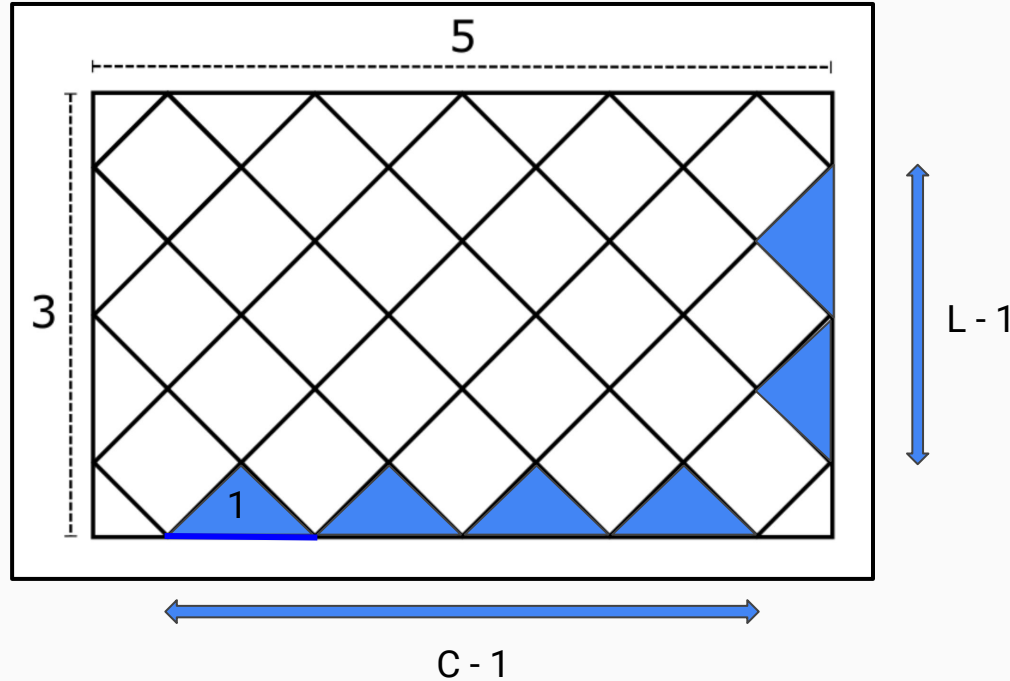
$$t1 = L * C + (L - 1)(C - 1)$$

$$t2 =$$



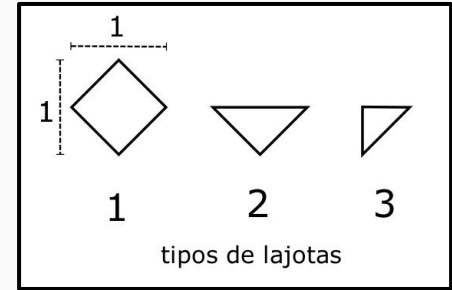
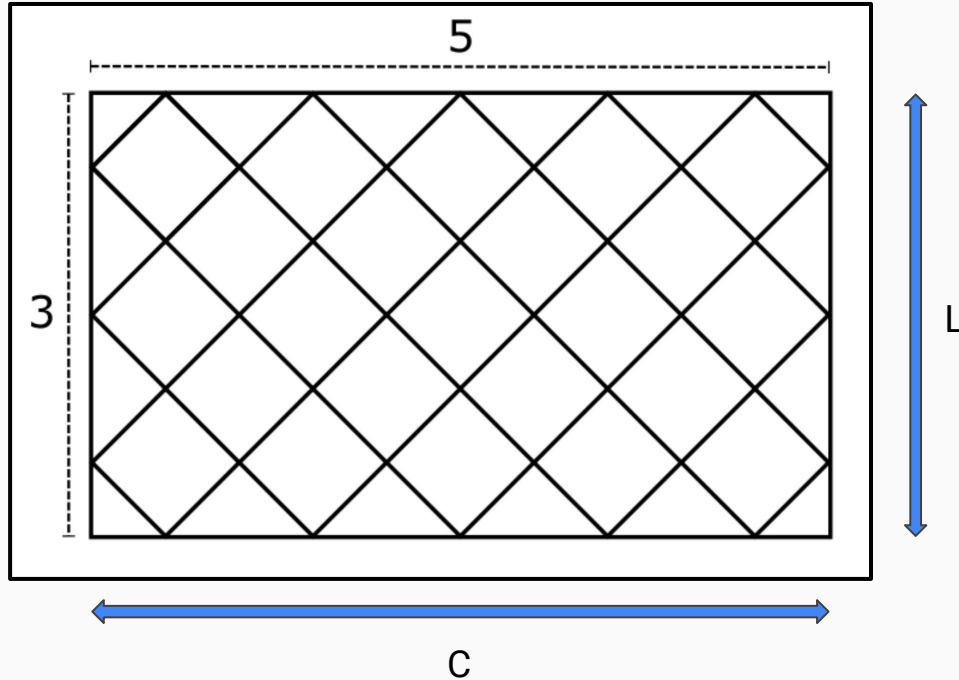
$$t1 = L * C + (L - 1)(C - 1)$$

$$t2 = (L - 1) + (C - 1)$$



$$t1 = L * C + (L - 1)(C - 1)$$

$$t2 = 2(L - 1) + 2(C - 1)$$



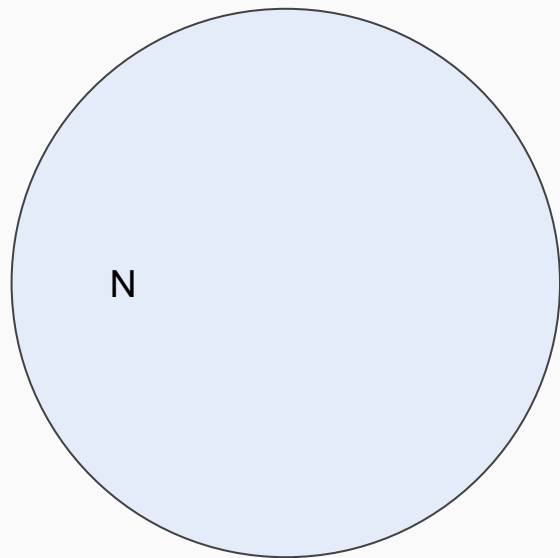
$$t1 = L * C + (L - 1)(C - 1)$$

$$t2 = 2(L - 1) + 2(C - 1)$$

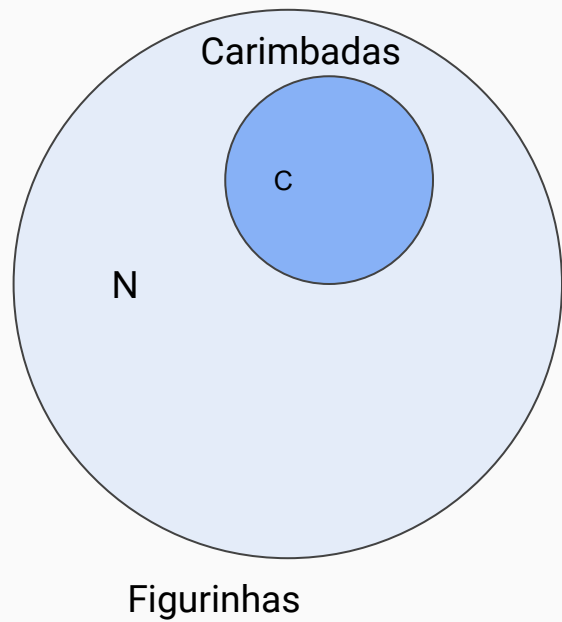
Implementação

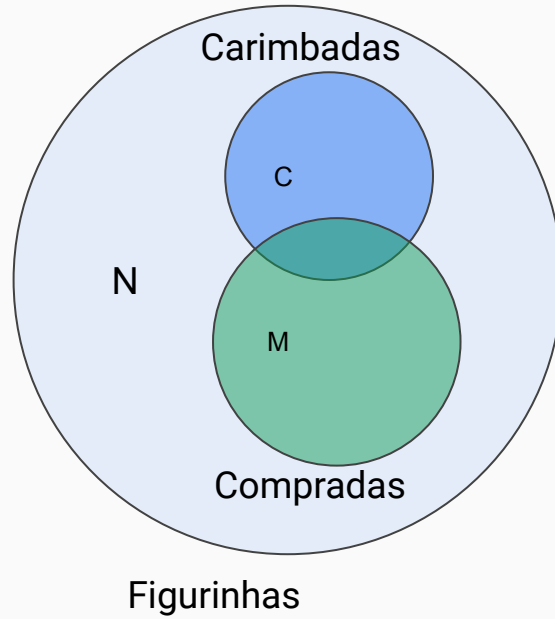
Figurinhas da copa

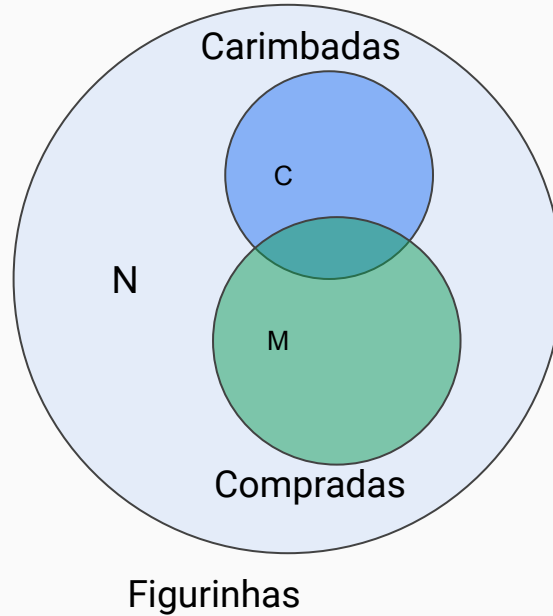
- Médio
- Estrutura de dado simples
- Laço de repetição



Figurinhas







Figurinhas carimbadas que não foram compradas

$$R = C - M$$

$$\{x \mid x \in C \text{ e } x \notin M\}$$

Entrada

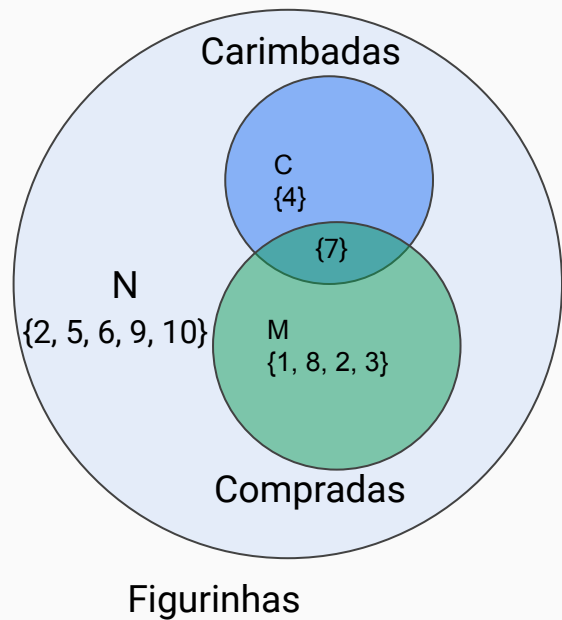
10 2 5

4 7

7 1 2 8 3

Saída

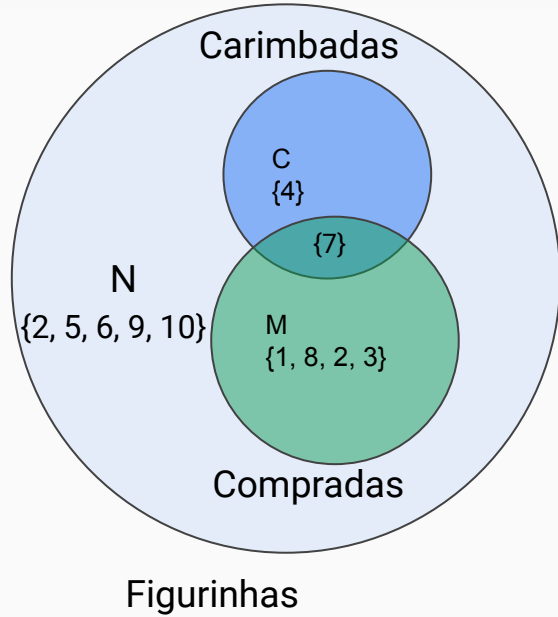
1



Figurinhas carimbadas que não foram compradas

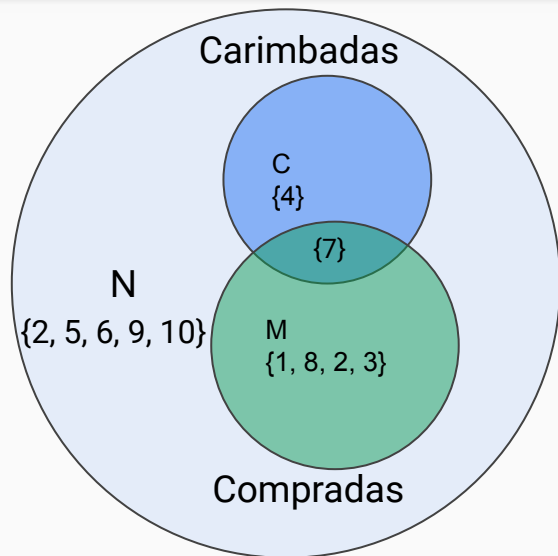
$$R = C - M$$

$$\{x \mid x \in C \text{ e } x \notin M\}$$



0	1	2	3	4	5	6	7	8	9	10

$N + 1$ posições



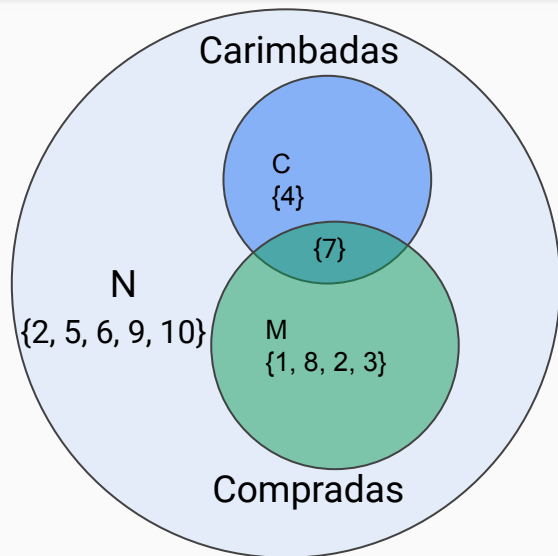
Figurinhas

0	1	2	3	4	5	6	7	8	9	10
0	0	0	0	0	0	0	0	0	0	0

$N + 1$ posições

Carimbadas = {4, 7}

Compradas = {7, 1, 8, 2, 3}



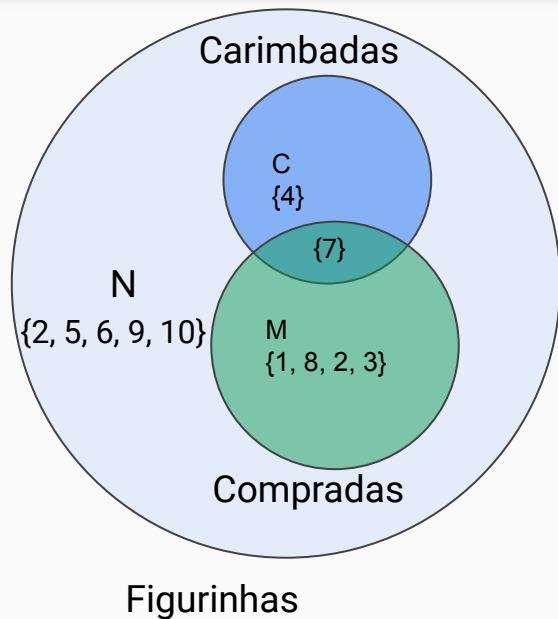
0	1	2	3	4	5	6	7	8	9	10
0	1	1	1	0	0	0	1	1	0	0

N + 1 posições

Carimbadas = {4, 7}

Compradas = {7, 1, 8, 2, 3}

Figurinhas



0	1	2	3	4	5	6	7	8	9	10
0	1	1	1	0	0	0	1	1	0	0

N + 1 posições

Carimbadas = $\{4, 7\}$

Compradas = $\{7, 1, 8, 2, 3\}$

Resposta = Quais figurinhas do conjunto de carimbadas não foram compradas

Índices que possuem valor 0 no vetor

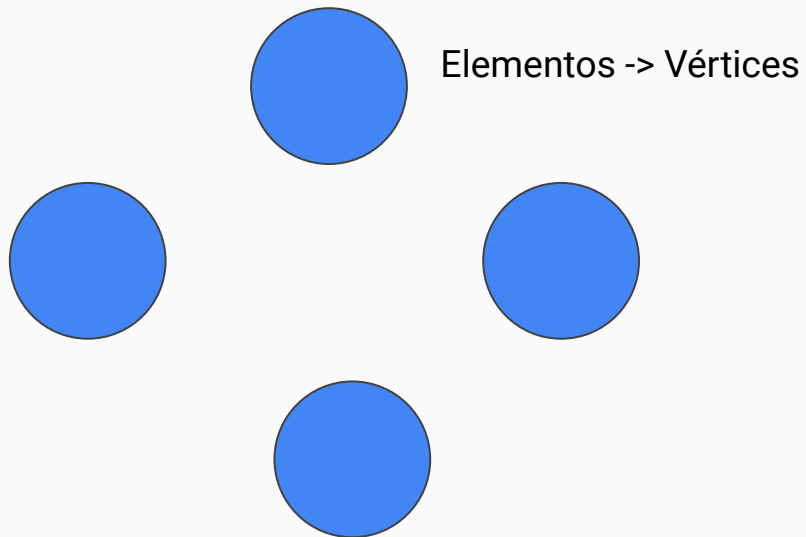
Implementação

Ilhas

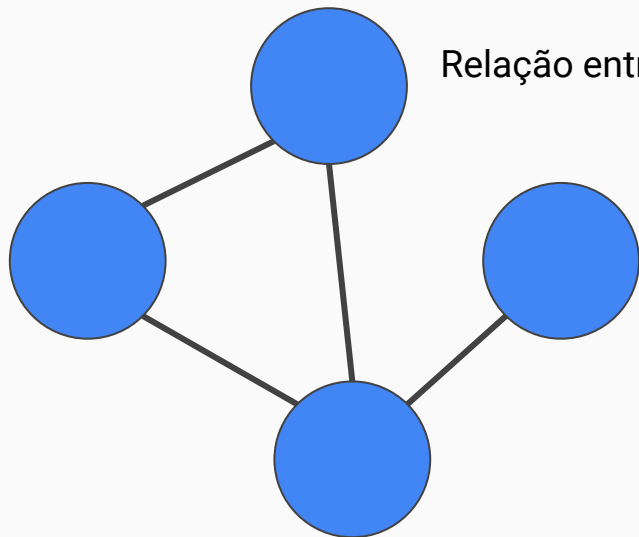
- Difícil
- Grafos
- Melhor caminho entre vértices
- Algoritmo de Dijkstra

O que são grafos?

O que são grafos?

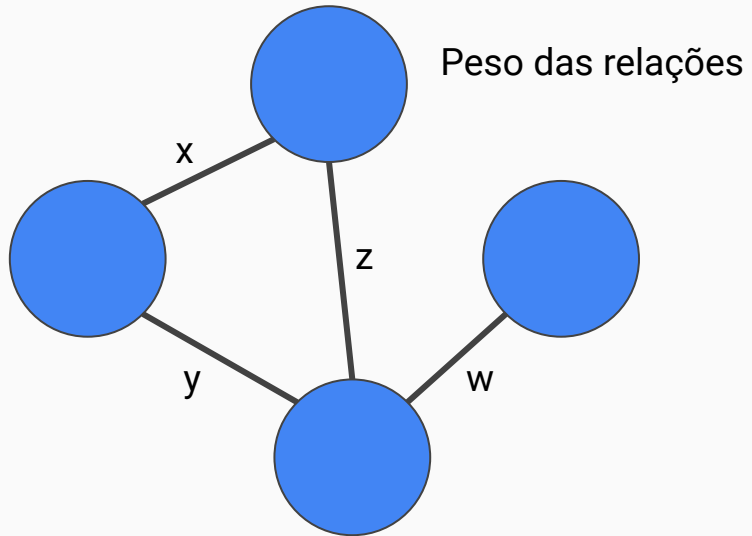


O que são grafos?



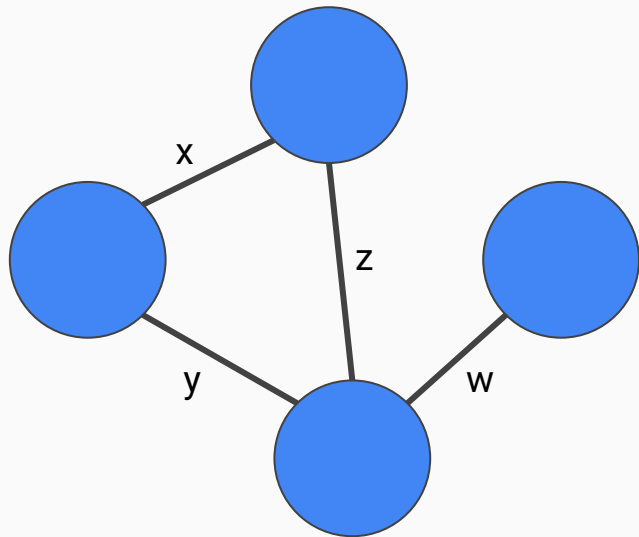
Relação entre elementos -> Aresta

O que são grafos?



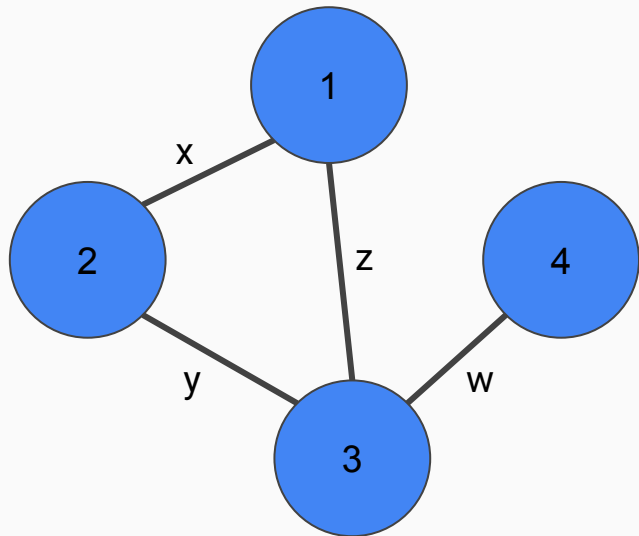
O que são grafos?

Como representar um grafo?

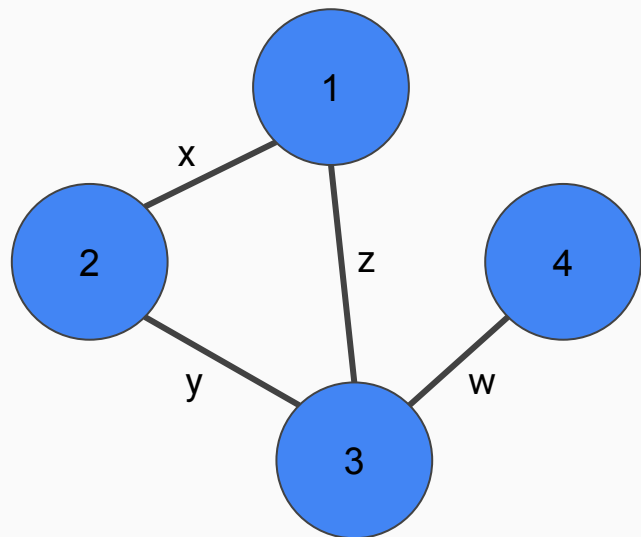


O que são grafos?

Como representar um grafo?



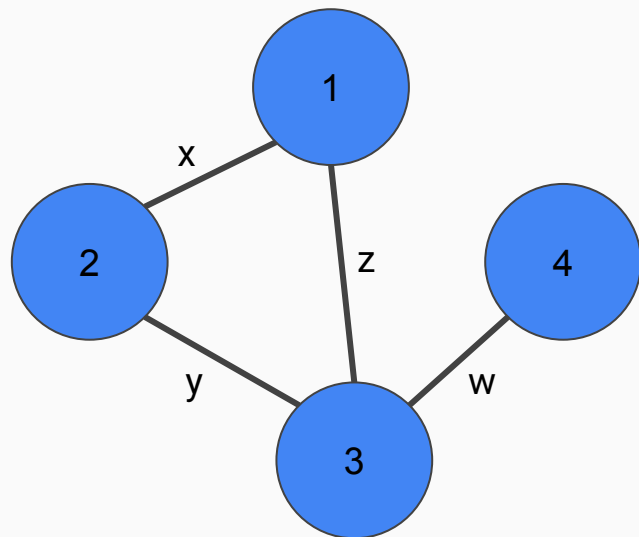
O que são grafos?



Como representar um grafo?

1	2 x	3 z
---	--------	--------

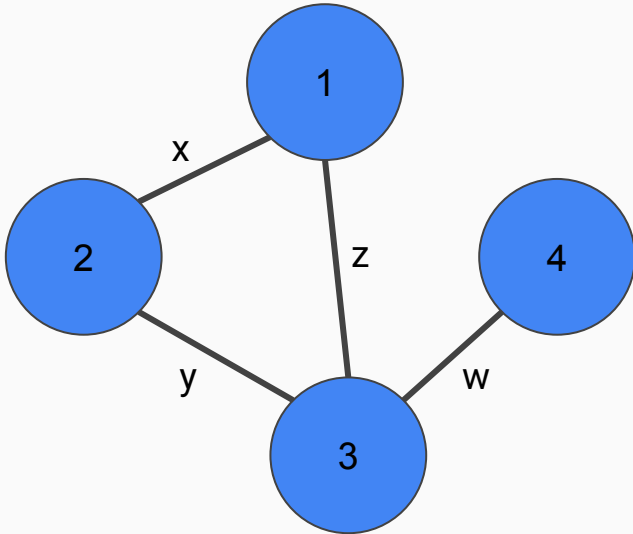
O que são grafos?



Como representar um grafo?

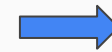
1	2 x	3 z	
2	1 x	3 y	
3	1 z	2 y	4 w
4	3 w		

O que são grafos?

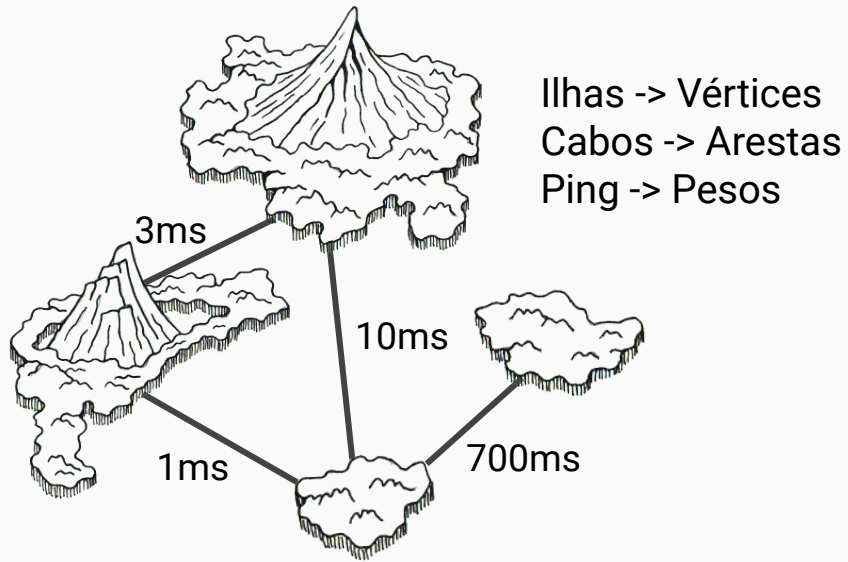


Como representar um grafo?

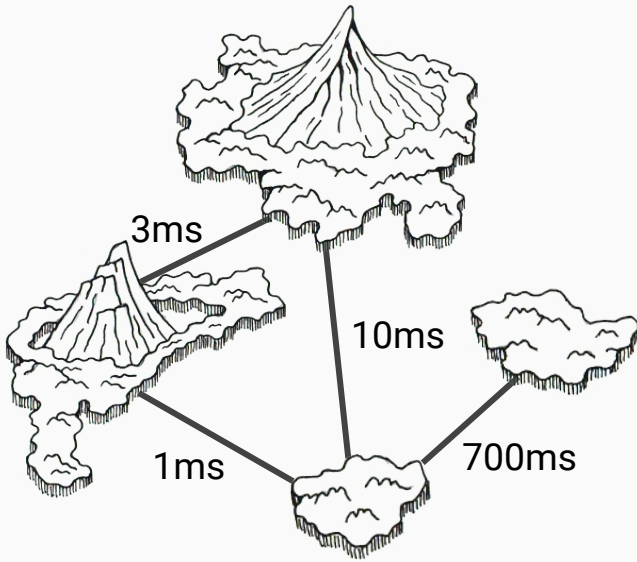
1	2 x	3 z	
2	1 x	3 y	
3	1 z	2 y	4 w
4	3 w		

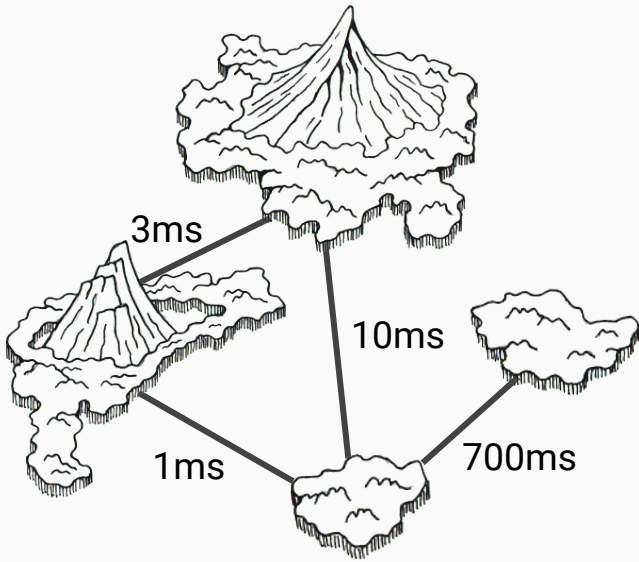


Lista de Adjacência



Ilha principal -> Servidor

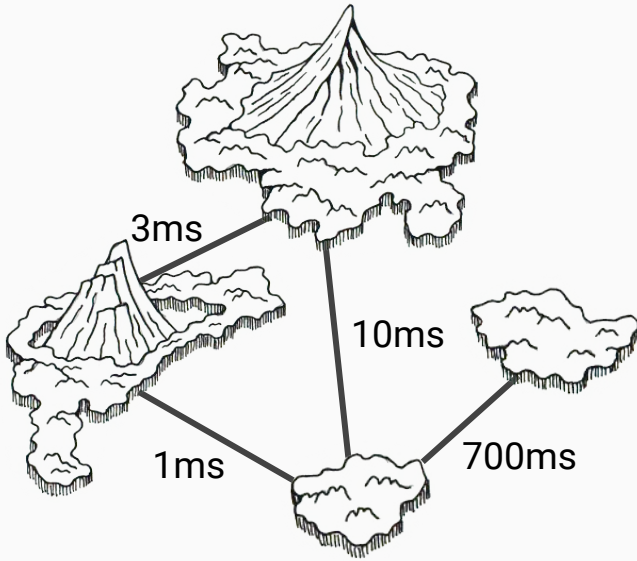




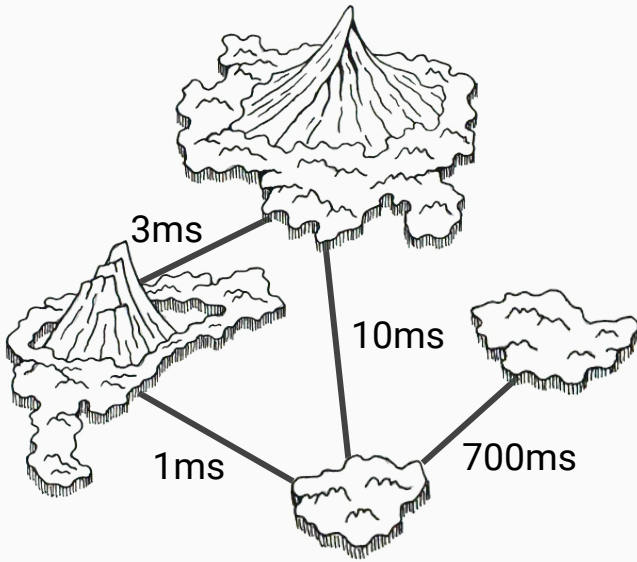
Ilha principal -> Servidor

Quando duas ilhas se comunicam através de uma série de cabos, o ping entre elas é a soma dos pings de cada cabo no caminho

Um par de ilhas sempre se comunica através do caminho com menor ping possível



Diferença dos pings entre a ilha com o maior ping e a ilha com o menor ping até o servidor.



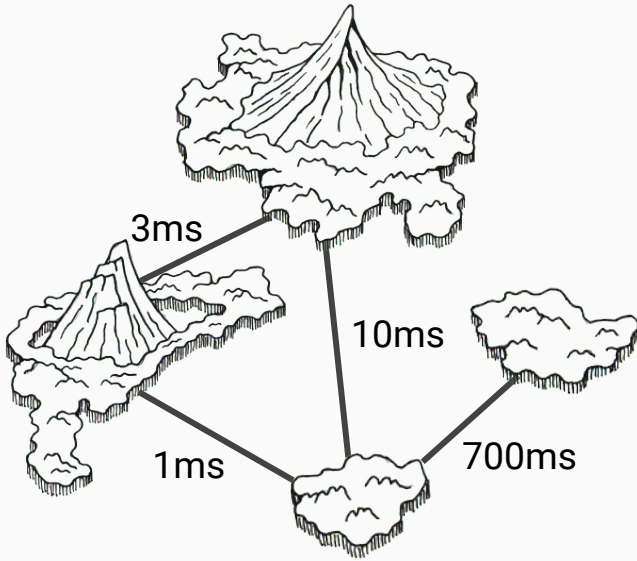
Diferença dos pings entre a ilha com o maior ping e a ilha com o menor ping até o servidor.

3

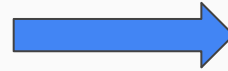
4

704

Maior - Menor = $704 - 3 = 701$



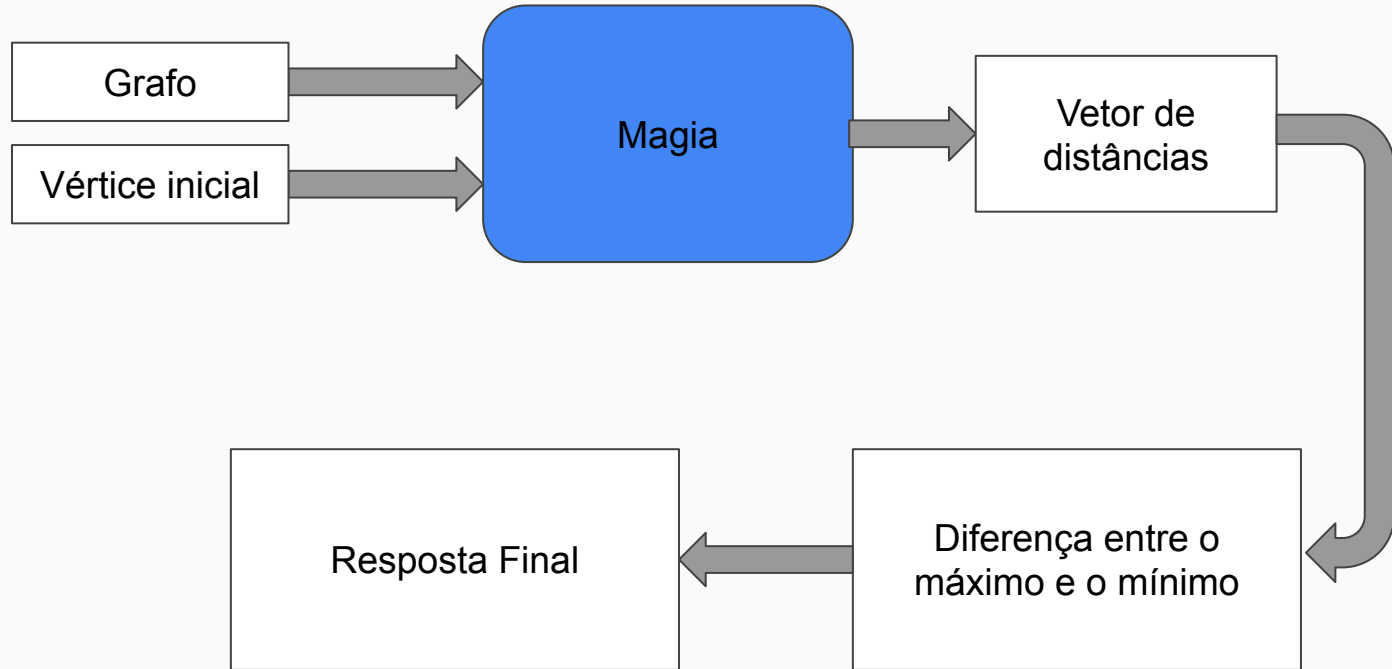
3
4
704

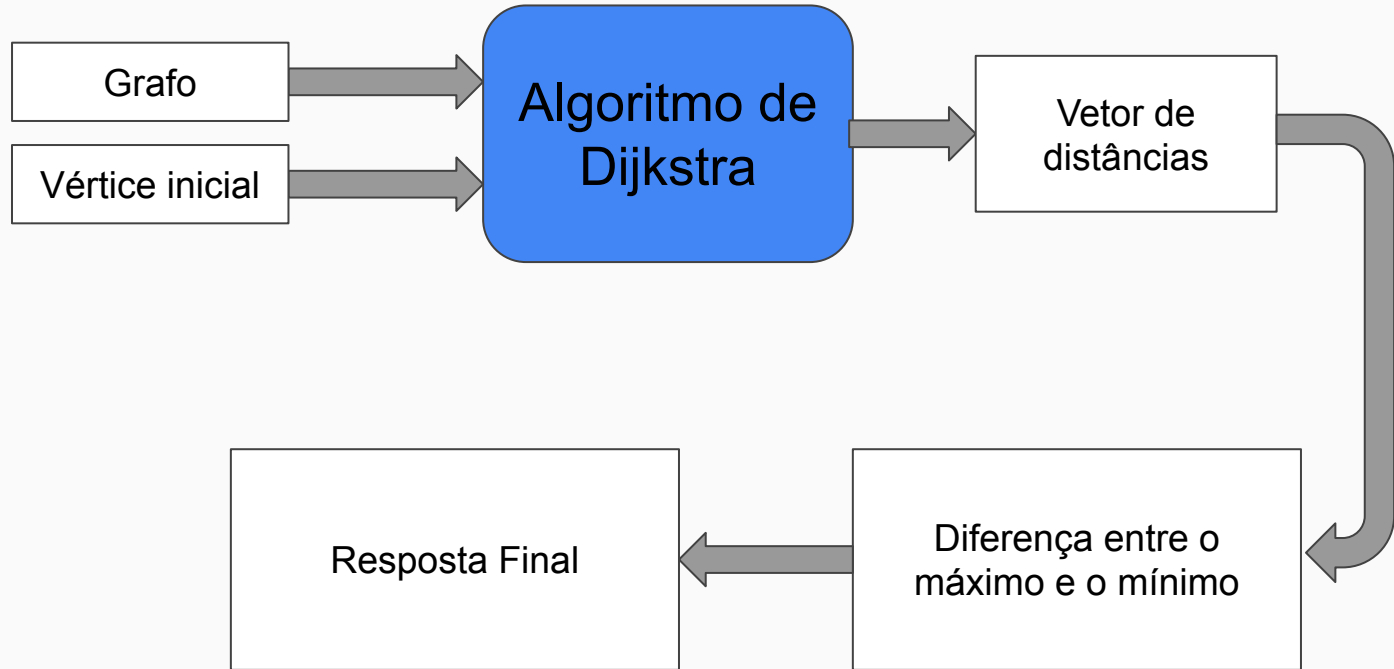


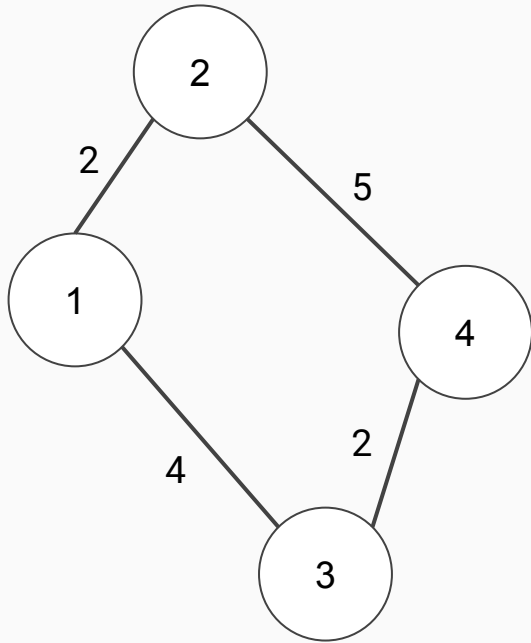
Vetor de
distâncias

Pensar em distâncias é mais intuitivo
São problemas análogos



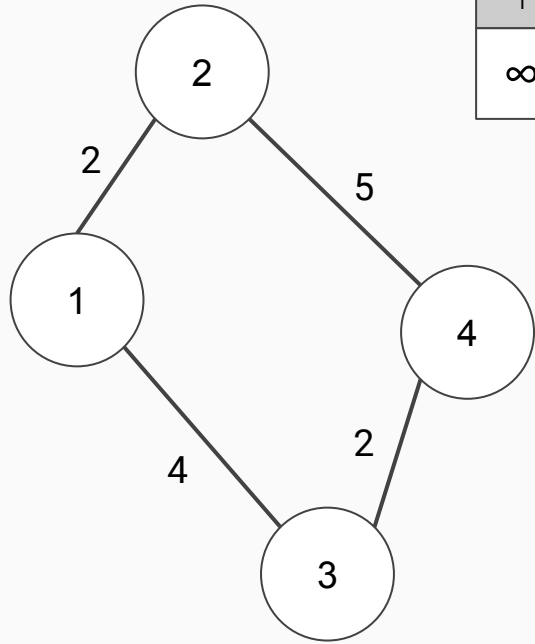




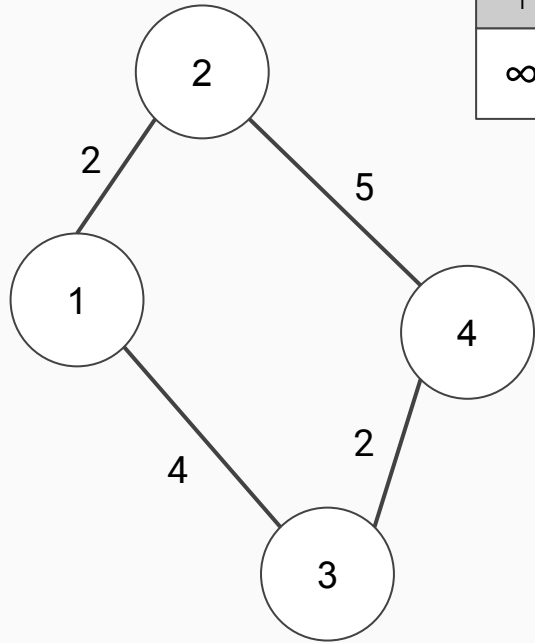


Servidor -> Vértice 1

Calcular o vetor de distâncias



1	2	3	4
∞	∞	∞	∞



1	2	3	4
∞	∞	∞	∞

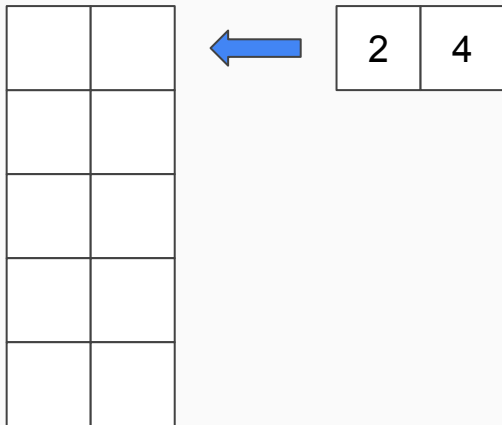
Fila de prioridade:

d	v

Fila de prioridade

Na fila, os menores elementos
aparecem nas primeiras posições

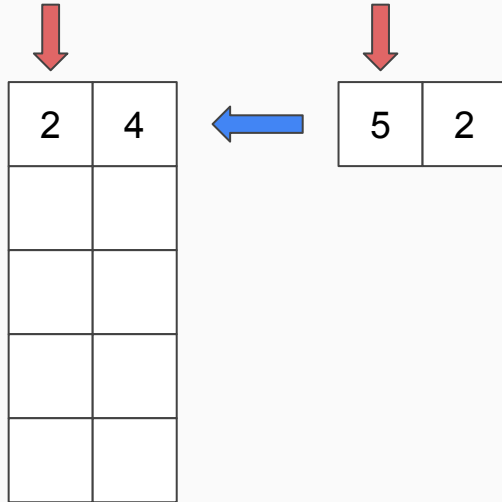
Fila de prioridade



Fila de prioridade

2	4

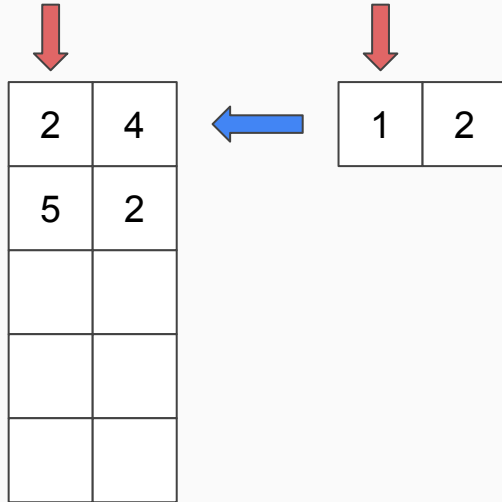
Fila de prioridade



Fila de prioridade

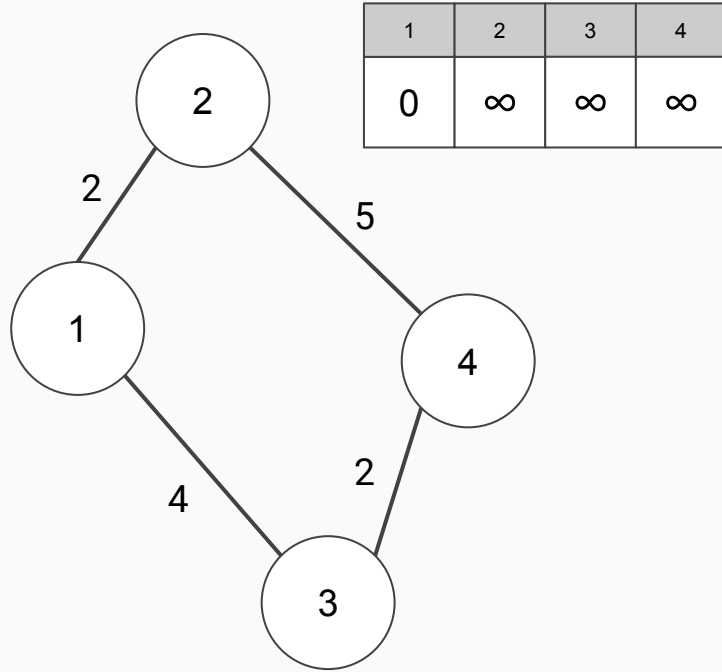
2	4
5	2

Fila de prioridade

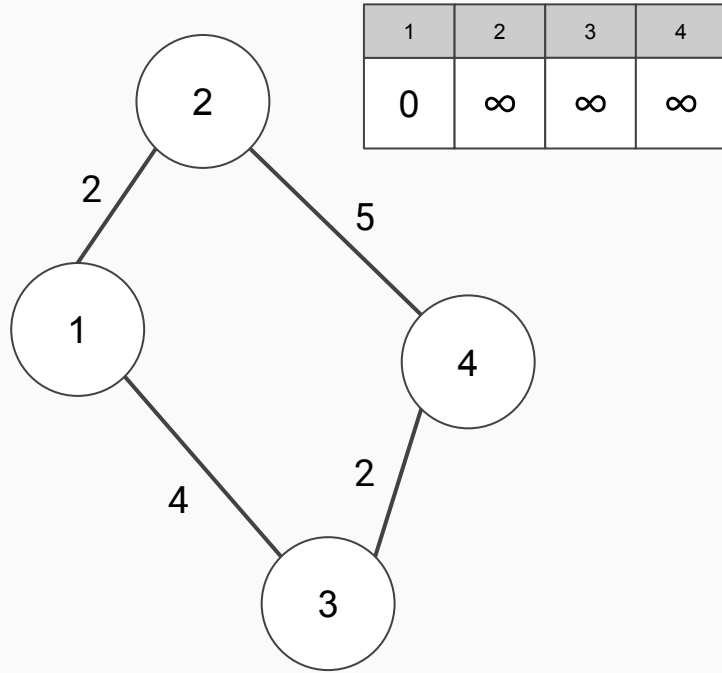


Fila de prioridade

1	2
2	4
5	2



d	v
0	1

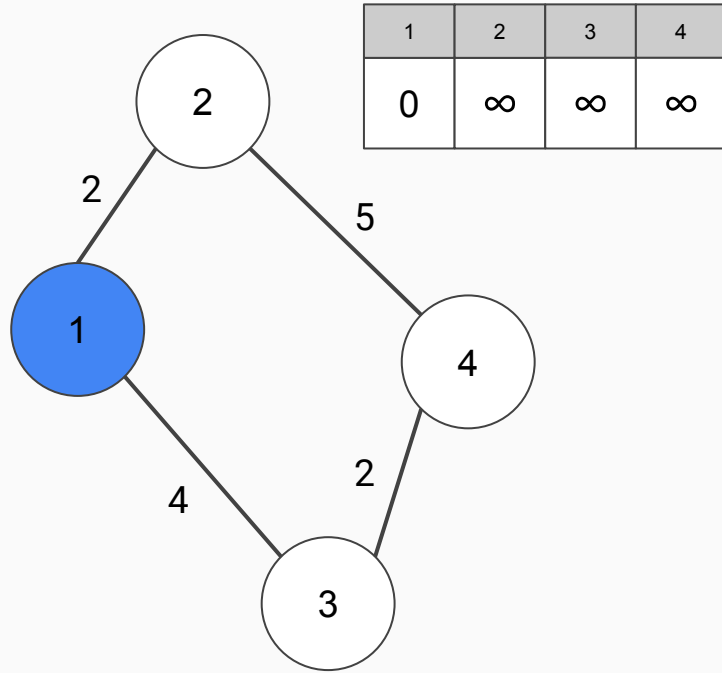


1	2	3	4
0	∞	∞	∞

0	1
---	---

1º Pegar o primeiro elemento da fila

d	v



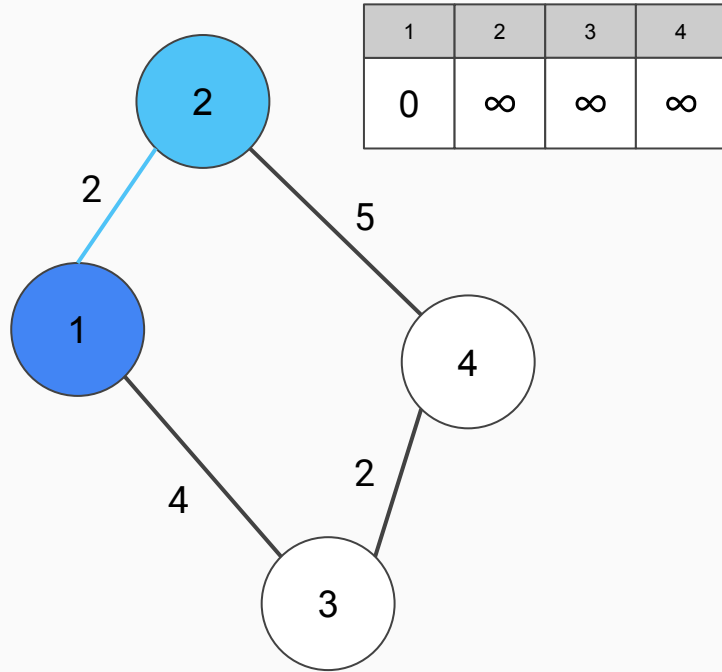
1	2	3	4
0	∞	∞	∞

0	1
---	---

2º Verificar se ele já foi visitado.

Se sim, não faça nada.
Se não, marcá-lo como visitado

d	v



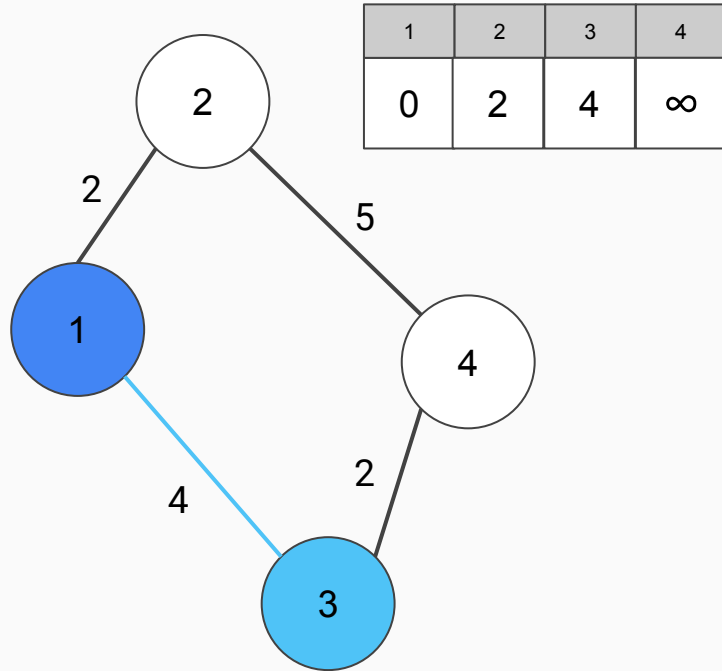
0	1
---	---

3º Processar esse vértice olhando para os vértices adjacentes a ele

2	2
---	---

Se $d[1] + \text{peso} < d[2]$:
 $d[2] = d[1] + \text{peso}$
 inclui o par $\{d[2], 2\}$ na fila

d	v



0	1
---	---

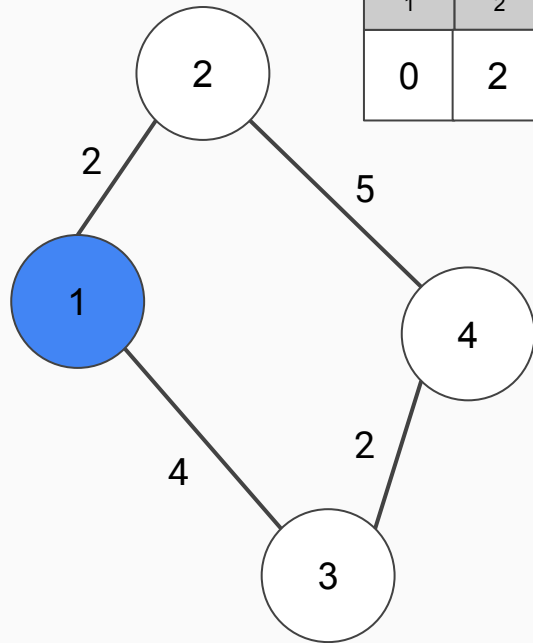
3º Processar esse vértice olhando para os vértices adjacentes a ele

4	3
---	---

Se $d[1] + \text{peso} < d[3]$:
 $d[3] = d[1] + \text{peso}$
 inclui o par $\{d[3], 3\}$ na fila

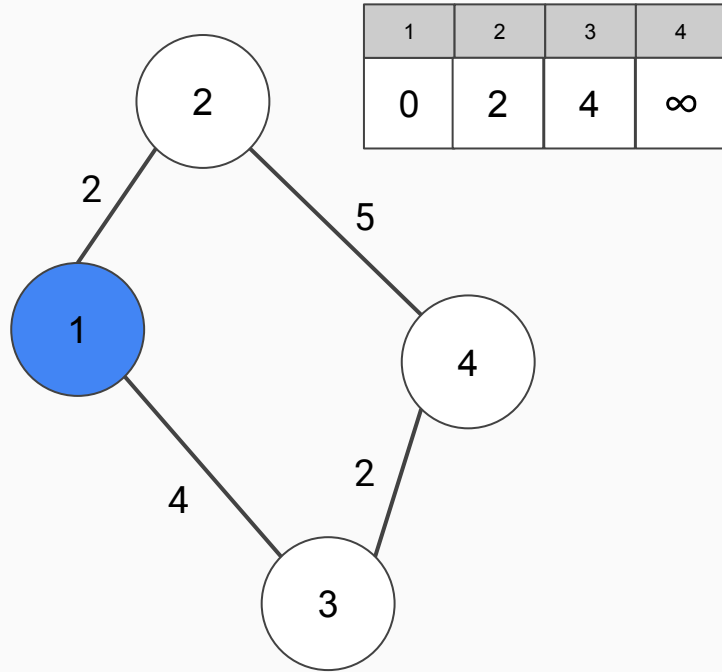
d	v
2	2
4	3

1	2	3	4
0	2	4	∞



4º Repita os passos 1 - 3 até a fila ficar vazia

d	v
2	2
4	3

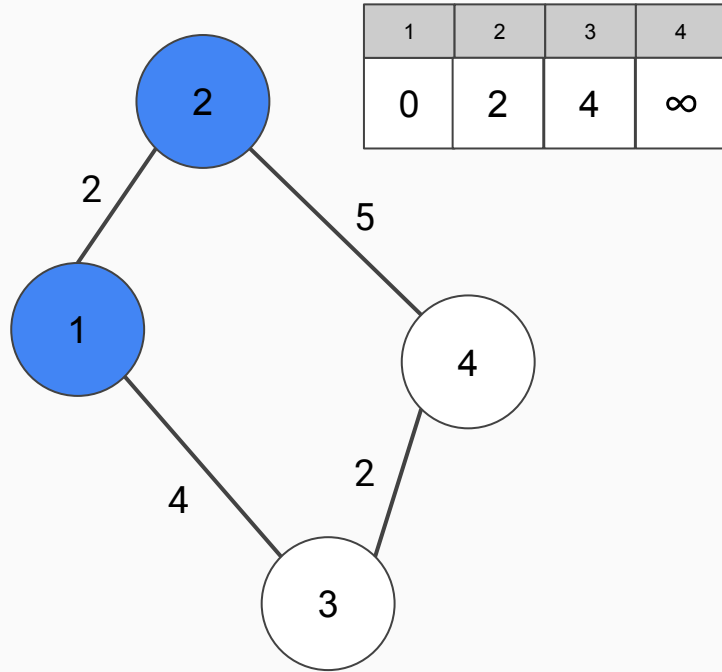


1	2	3	4
0	2	4	∞

2	2
---	---

1º Pegar o primeiro elemento da fila

d	v
4	3

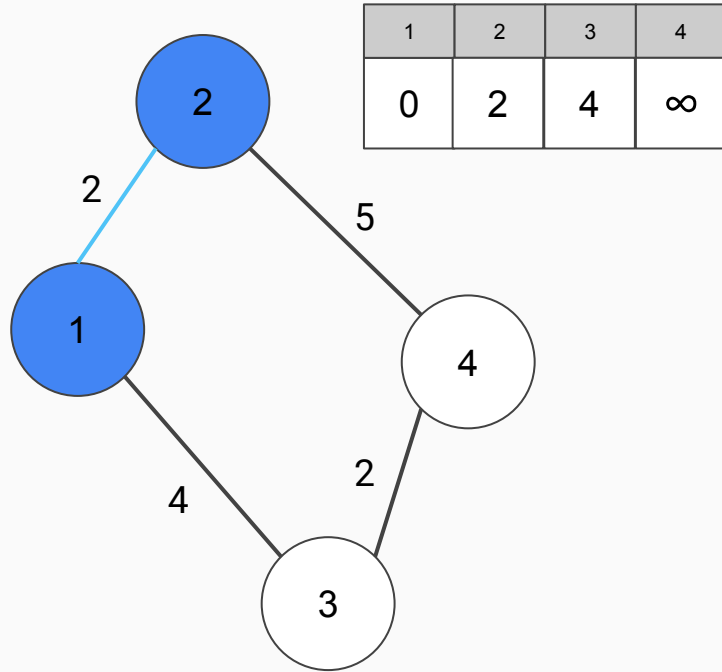


2	2
---	---

2º Verificar se ele já foi visitado.

Se sim, não faça nada.
Se não, marcá-lo como visitado

d	v
4	3



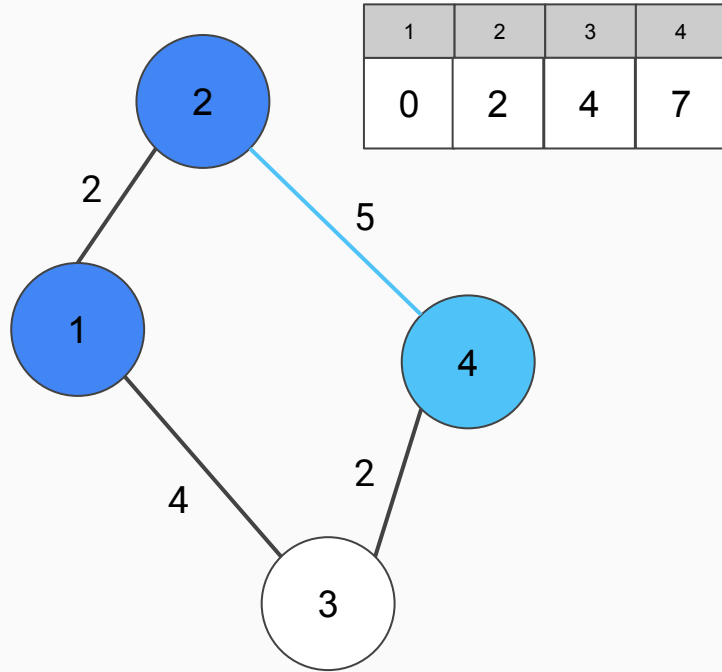
2	2
---	---

3º Processar esse vértice olhando para os vértices adjacentes a ele

2	1
---	---

Se $d[2] + \text{peso} < d[1]$:
 $d[1] = d[2] + \text{peso}$
 inclui o par $\{d[1], 1\}$ na fila

d	v
4	3



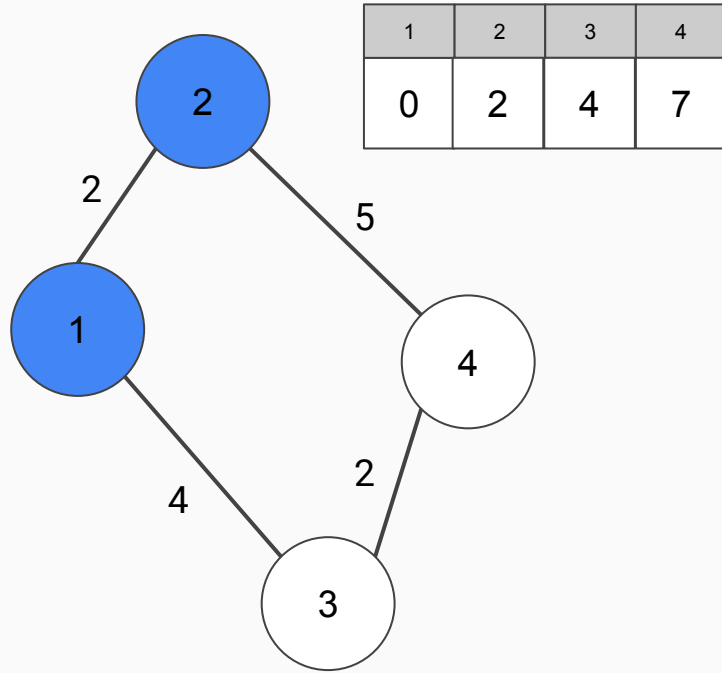
2	2
---	---

3º Processar esse vértice olhando para os vértices adjacentes a ele

5	4
---	---

Se $d[2] + \text{peso} < d[4]$:
 $d[4] = d[2] + \text{peso}$
 inclui o par $\{d[4], 4\}$ na fila

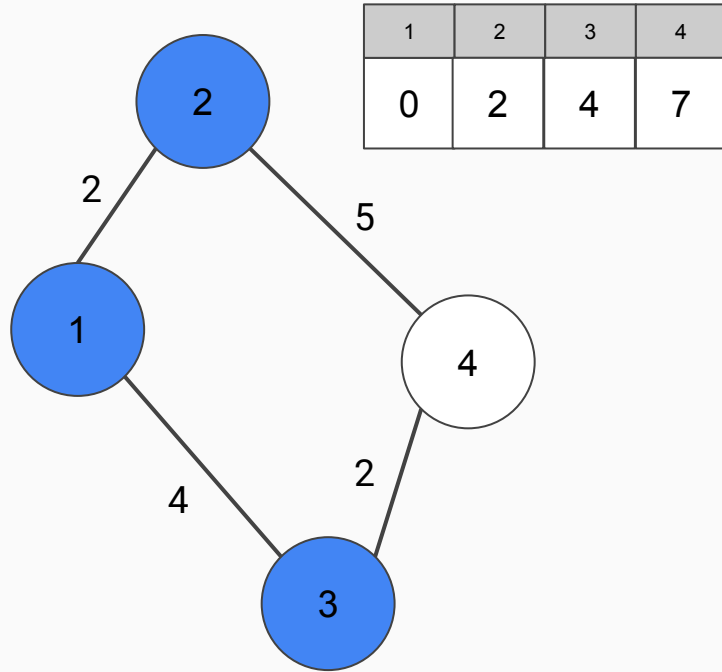
d	v
4	3
7	4



4	3
---	---

1º Pegar o primeiro elemento da fila

d	v
7	4

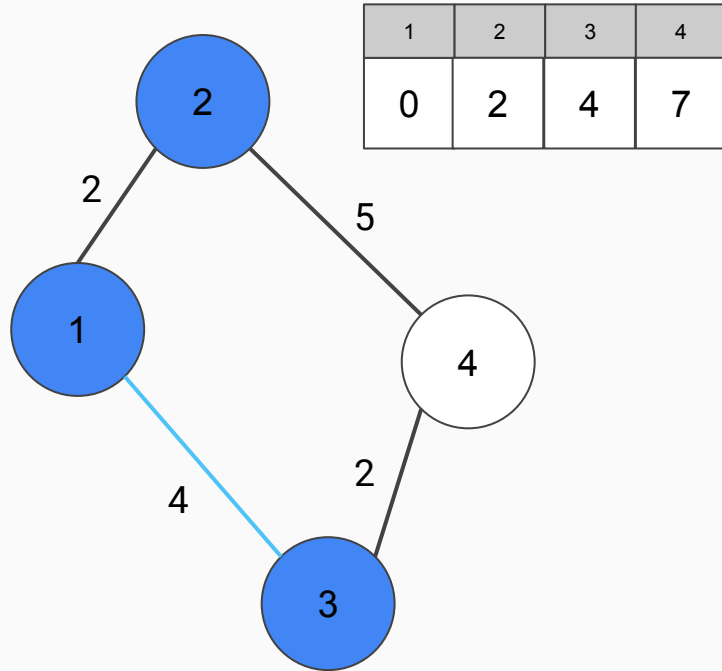


4	3
---	---

2º Verificar se ele já foi visitado.

Se sim, não faça nada.
Se não, marcá-lo como visitado

d	v
7	4



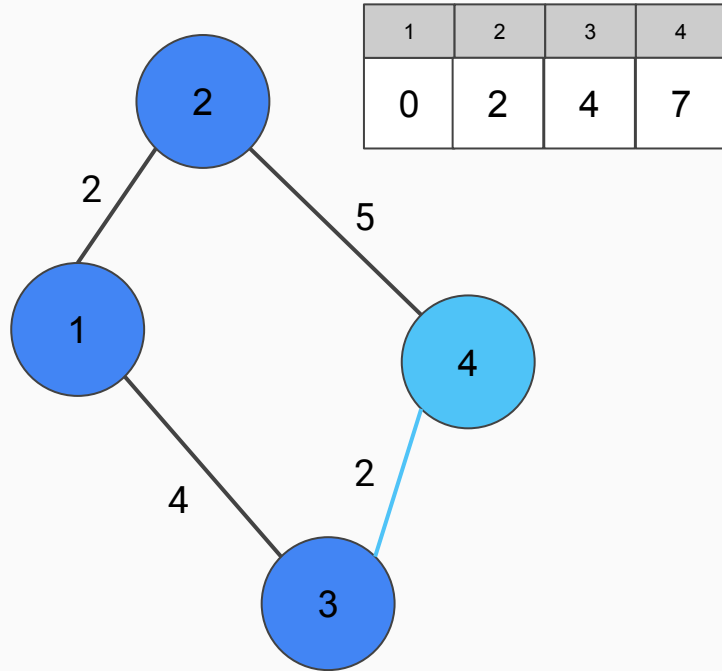
4	3
---	---

3º Processar esse vértice olhando para os vértices adjacentes a ele

4	1
---	---

Se $d[1] + \text{peso} < d[3]$:
 $d[1] = d[3] + \text{peso}$
 inclui o par $\{d[1], 1\}$ na fila

d	v
7	4



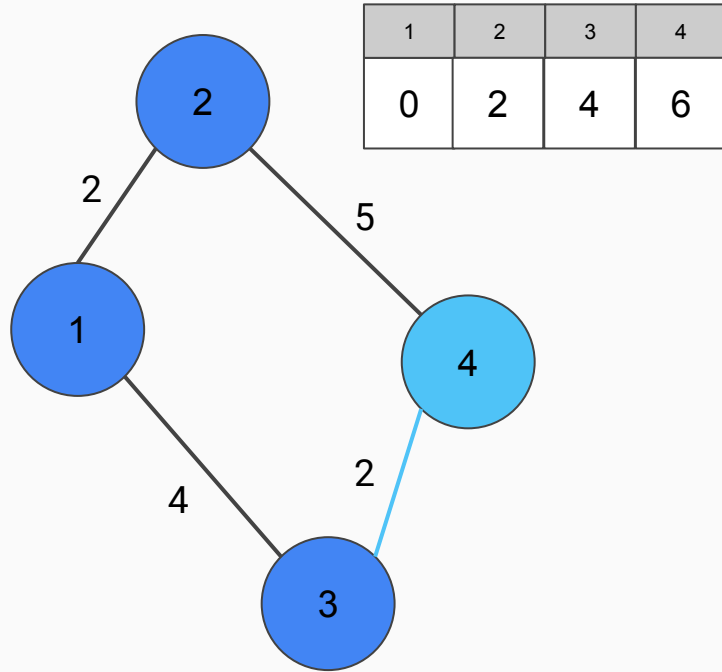
4	3
---	---

3º Processar esse vértice olhando para os vértices adjacentes a ele

2	4
---	---

Se $d[4] + \text{peso} < d[3]$:
 $d[4] = d[3] + \text{peso}$
 inclui o par $\{d[4], 4\}$ na fila

d	v
7	4



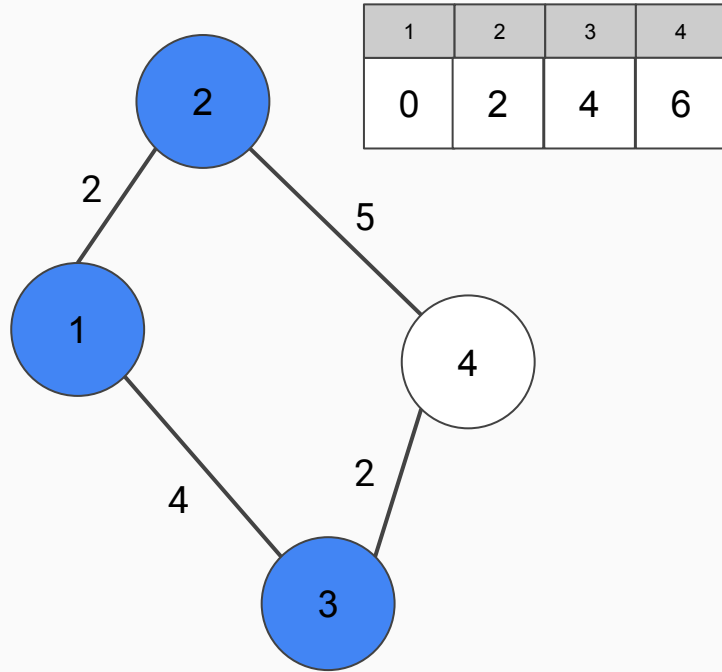
4	3
---	---

3º Processar esse vértice olhando para os vértices adjacentes a ele

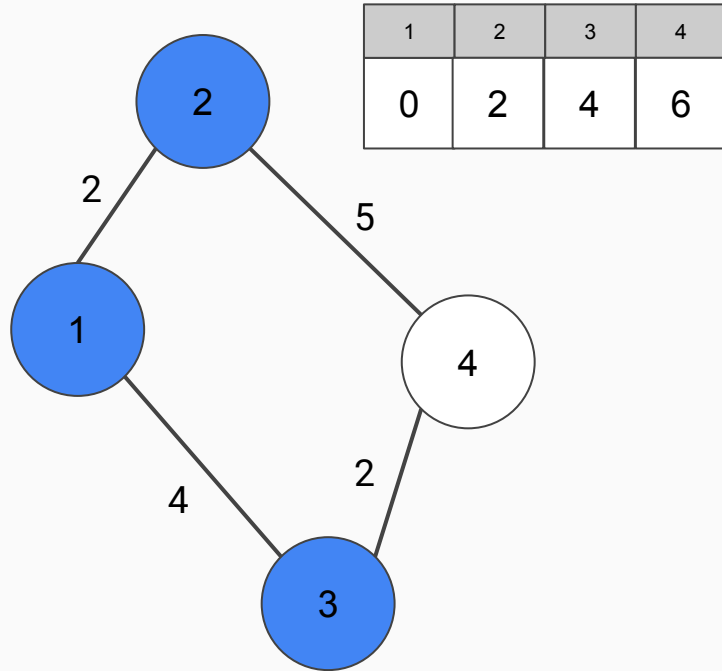
2	4
---	---

Se $d[4] + \text{peso} < d[3]$:
 $d[4] = d[3] + \text{peso}$
 inclui o par $\{d[4], 4\}$ na fila

d	v
6	4
7	4

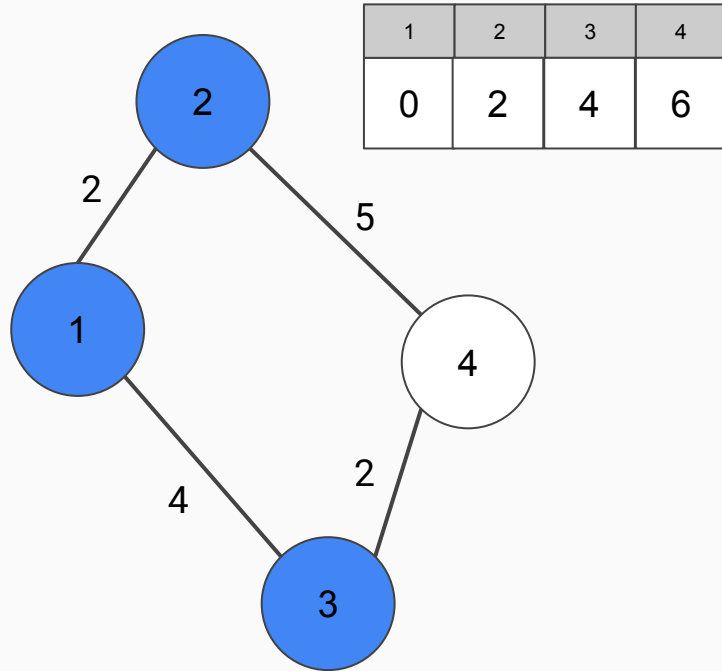


d	v
6	4
7	4



Existem dois caminhos do 1 até o 4, um que demora 7, outro que demora 6

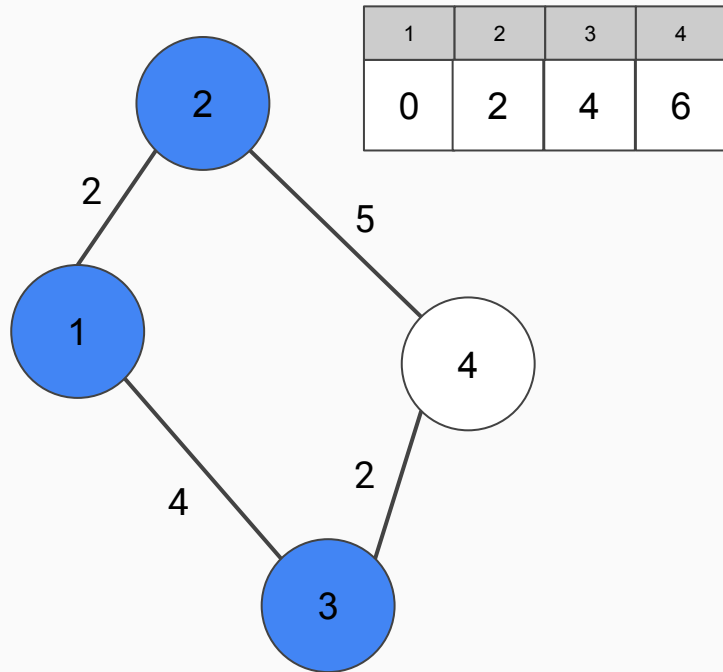
d	v
6	4
7	4



Existem dois caminhos do 1 até o 4, um que demora 7, outro que demora 6

d	v
6	4
7	4

A fila de prioridade faz com que eu sempre olhe as menores distâncias primeiro

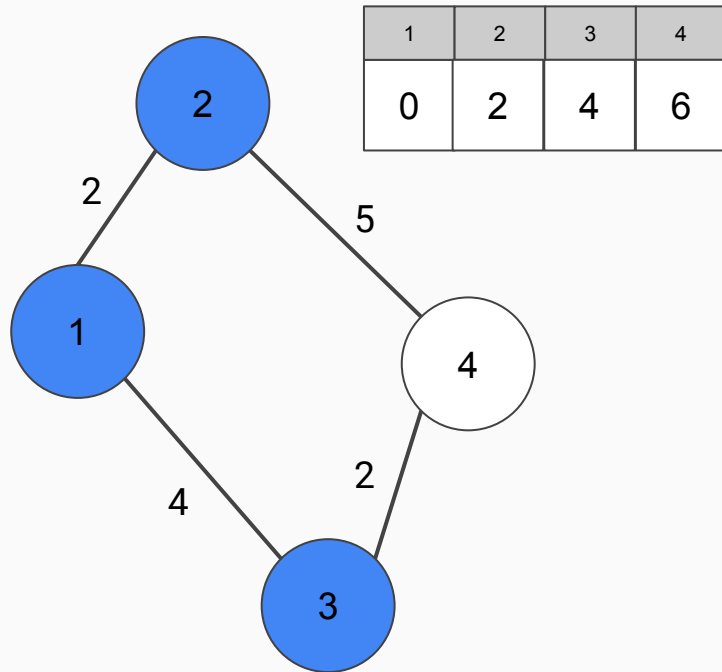


d	v
6	4
7	4

Existem dois caminhos do 1 até o 4, um que demora 7, outro que demora 6

A fila de prioridade faz com que eu sempre olhe as menores distâncias primeiro

Depois de processar o par {6, 4}, eu não irei olhar o par {7, 4}, pois eu já marquei que o 4 foi visitado



d	v
6	4
7	4

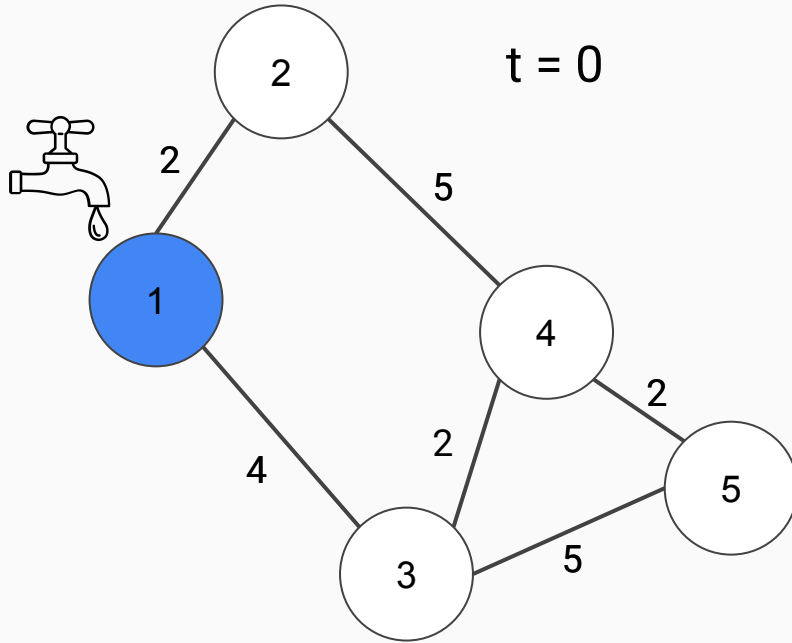
Existem dois caminhos do 1 até o 4, um que demora 7, outro que demora 6

A fila de prioridade faz com que eu sempre olhe as menores distâncias primeiro

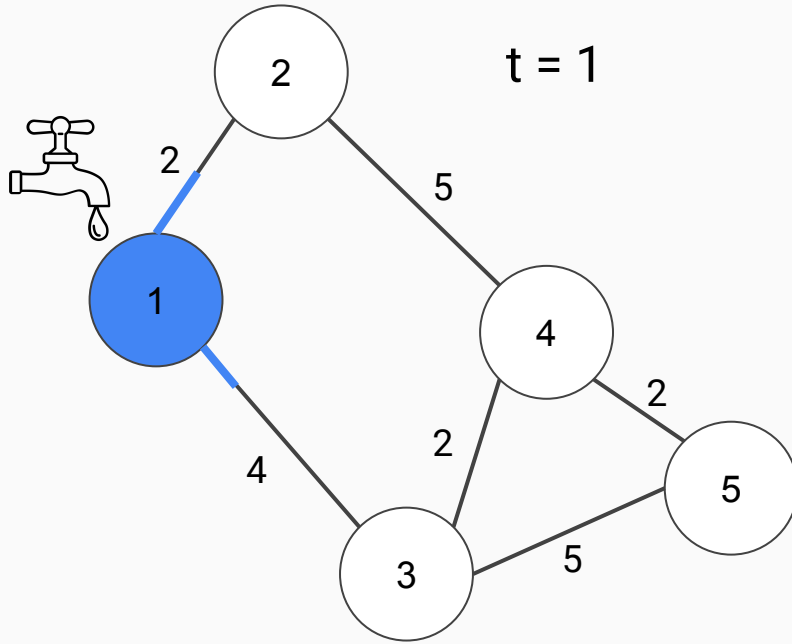
Depois de processar o par {6, 4}, eu não irei olhar o par {7, 4}, pois eu já marquei que o 4 foi visitado

Isso garante que eu sempre olhos os melhores caminhos até certo ponto

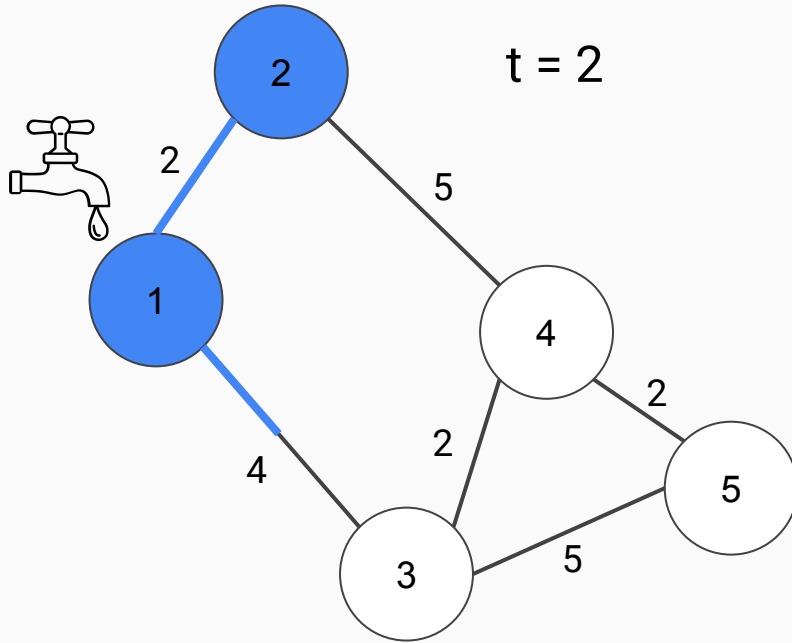
Porque funciona?



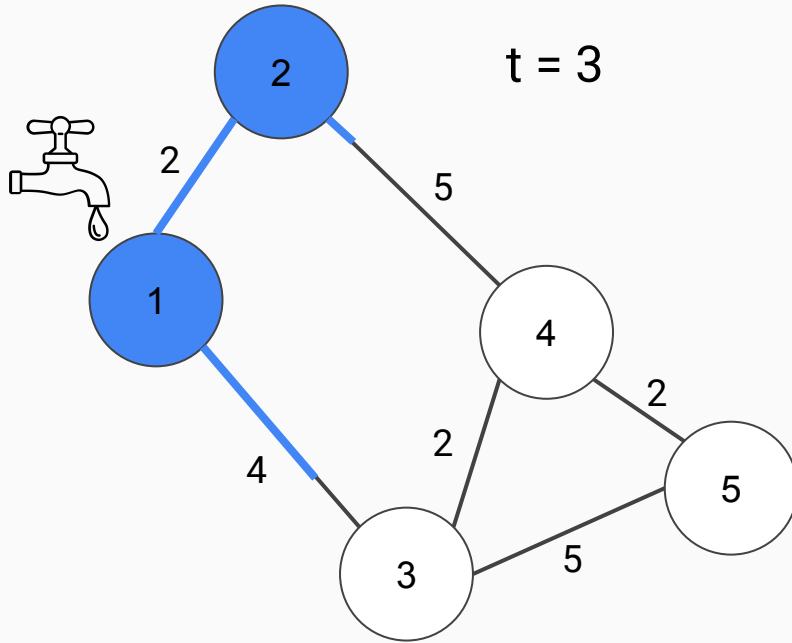
1	2	3	4	5
0				



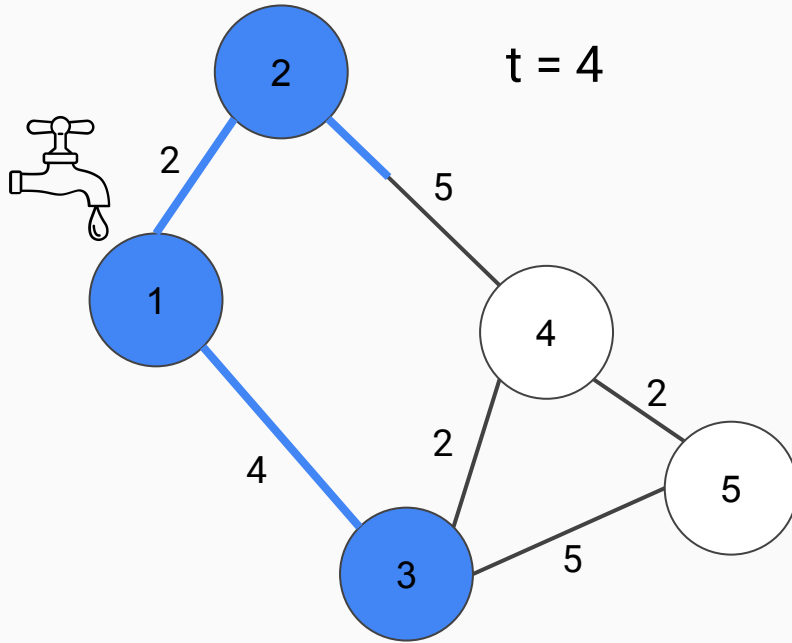
1	2	3	4	5
0				



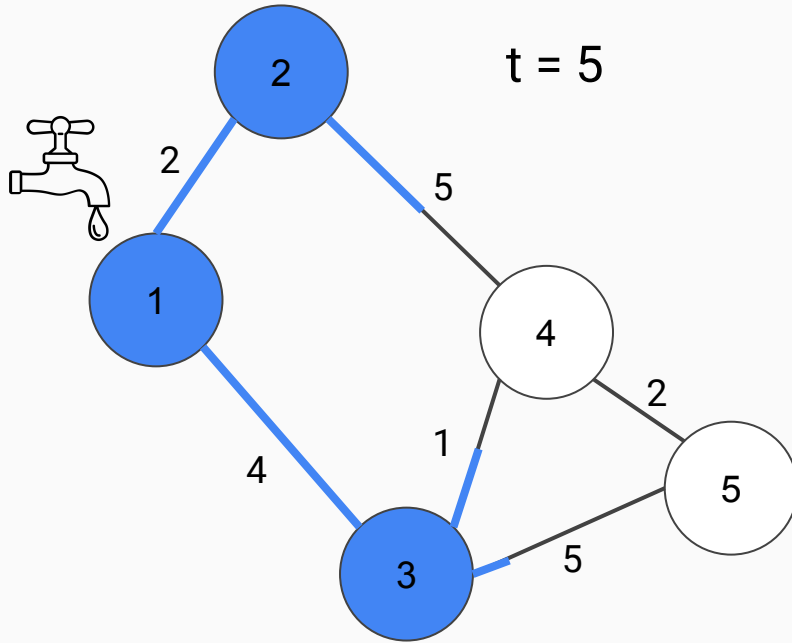
1	2	3	4	5
0	2			



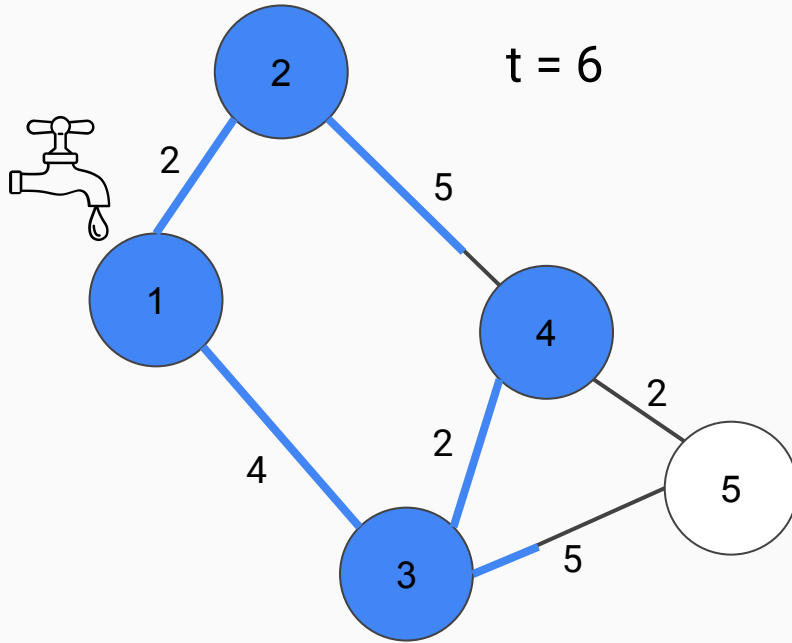
1	2	3	4	5
0	2			



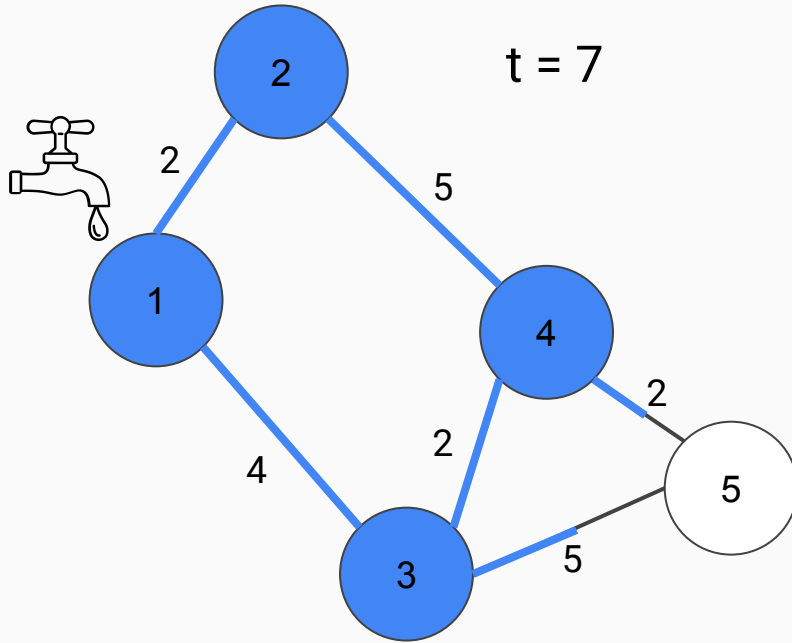
1	2	3	4	5
0	2	4		



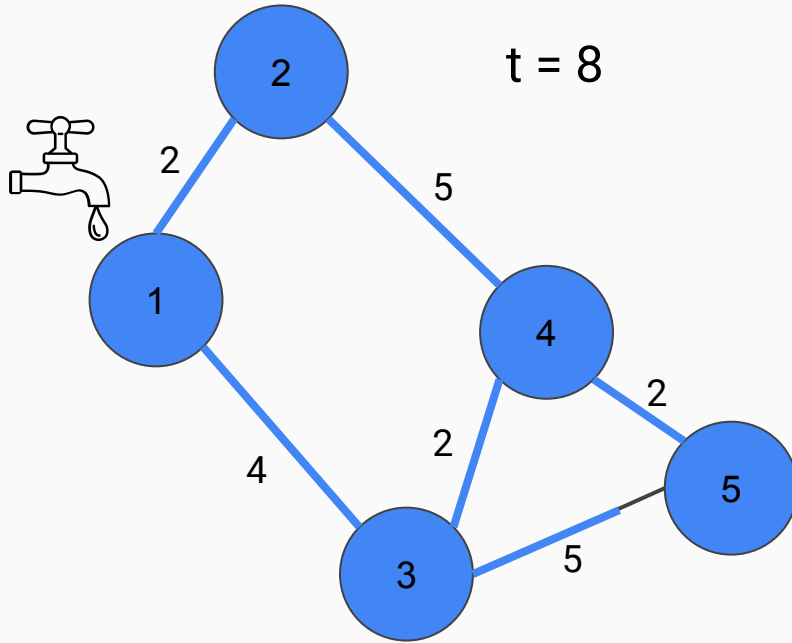
1	2	3	4	5
0	2	4		



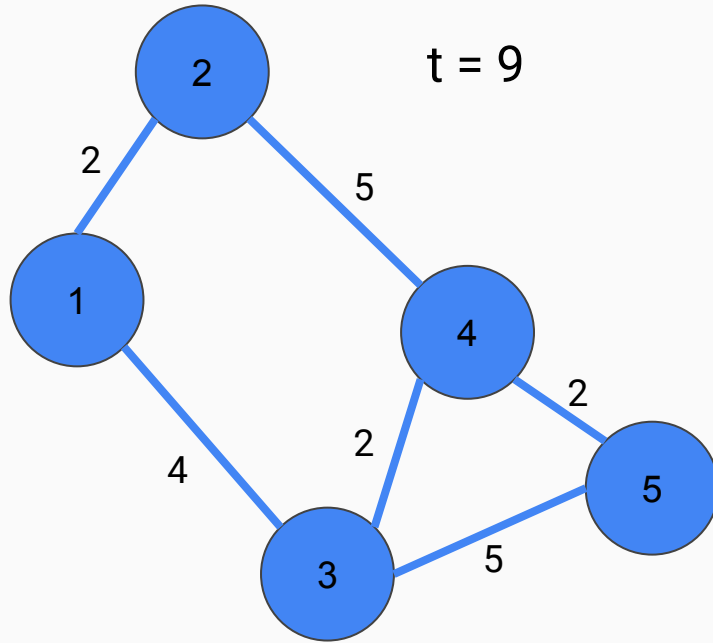
1	2	3	4	5
0	2	4	6	



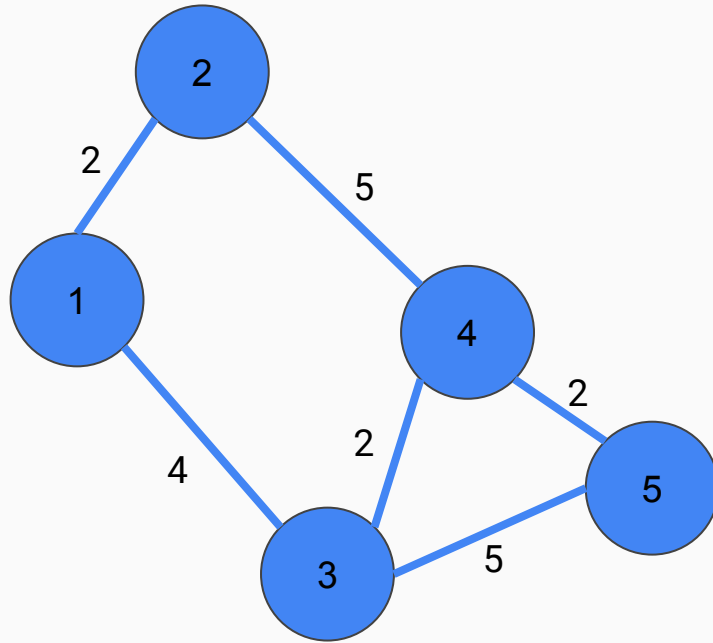
1	2	3	4	5
0	2	4	6	



1	2	3	4	5
0	2	4	6	8



1	2	3	4	5
0	2	4	6	8



1	2	3	4	5
0	2	4	6	8



Vetor de distâncias

Implementação

Visualização do Algoritmo

[Pathfinding Visualizer](#)

