# CS631 - Advanced Programming in the UNIX Environment Interprocess Communication

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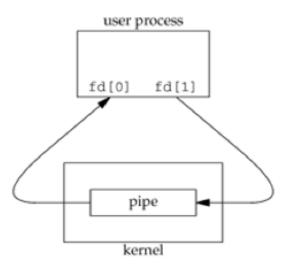
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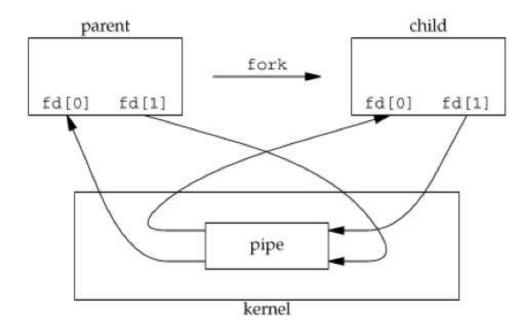
https://www.cs.stevens.edu/~jschauma/631/

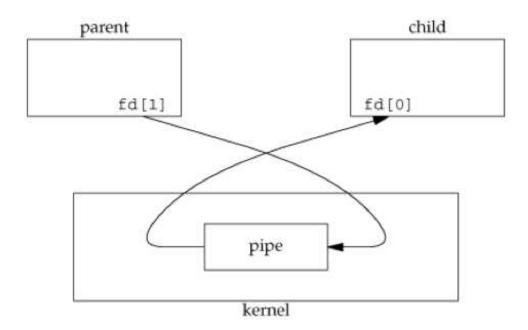
```
#include <unistd.h>
int pipe(int filedes[2]);

Returns: 0 if OK, -1 otherwise
```

- oldest and most common form of UNIX IPC
- half-duplex (on some versions full-duplex)







```
$ cc -Wall pipe1.c
$ ./a.out
P=> Parent process with pid 23988 (and its ppid 7474).
P=> Sending a message to the child process (pid 23989):
C=> Child process with pid 23989 (and its ppid 23988).
C=> Reading a message from the parent (pid 23988):
Hello child! I'm your parent pid 23988!
$
```

```
$ cc -Wall pipe1.c
$ ./a.out
P=> Parent process with pid 23984 (and its ppid 7474).
P=> Sending a message to the child process (pid 23985):
C=> Child process with pid 23985 (and its ppid 1).
C=> Reading a message from the parent (pid 1):
Hello child! I'm your parent pid 23984!
$
```

```
$ cc -Wall pipe1.c
$ ./a.out
P=> Parent process with pid 23986 (and its ppid 7474).
P=> Sending a message to the child process (pid 23987):
C=> Child process with pid 23987 (and its ppid 23986).
C=> Reading a message from the parent (pid 1):
Hello child! I'm your parent pid 23986!
$
```

A more useful example: displaying some content using the user's preferred pager. (Look, Ma, no temporary files!)

```
$ cat pipe2.c | ${PAGER:-/usr/bin/more}
$ cc -Wall pipe2.c
$ echo $PAGER
$ ./a.out pipe2.c
[...]
$ ps -o pid,ppid,command
  PTD PPTD COMMAND
22306 26650 ./a.out pipe2.c
22307 22306 more
23198 26650 ps -o pid, ppid, command
26650 26641 -ksh
$ fg
$ env PAGER=/bin/cat ./a.out pipe2.c
```

```
#include <unistd.h>
int pipe(int filedes[2]);

Returns: 0 if OK, -1 otherwise
```

- oldest and most common form of UNIX IPC
- half-duplex (on some versions full-duplex)
- can only be used between processes that have a common ancestor
- can have multiple readers/writers (PIPE\_BUF bytes are guaranteed to not be interleaved)

#### Behavior after closing one end:

- read(2) from a pipe whose write end has been closed returns 0 after all data has been read
- write(2) to a pipe whose read end has been closed generates SIGPIPE signal. If caught or ignored, write(2) returns an error and sets errno to EPIPE.

# Pipes: popen(3) and pclose(3)

- historically implemented using unidirectional pipe (nowadays frequently implemented using sockets or full-duplex pipes)
- type one of "r" or "w" (or "r+" for bi-directional communication, if available)
- cmd passed to /bin/sh -c

# Pipes: popen(3) and pclose(3)

```
$ cc -Wall popen.c
$ echo $PAGER

$ ./a.out popen.c
[...]
$ env PAGER=/bin/cat ./a.out popen.c
[...]
$
```

# Pipes: popen(3) and pclose(3)

```
$ cc -Wall popen.c
$ echo $PAGER

$ ./a.out popen.c
[...]
$ env PAGER=/bin/cat ./a.out popen.c
[...]
$ env PAGER=/bin/cat/foo ./a.out popen.c
sh: /bin/cat/foo: Not a directory
$ env PAGER="more; touch /tmp/boo" ./a.out popen.c
$ env PAGER="more; rm /etc/passwd 2>/dev/null" ./a.out popen.c
```

#### FIFOs: mkfifo(2)

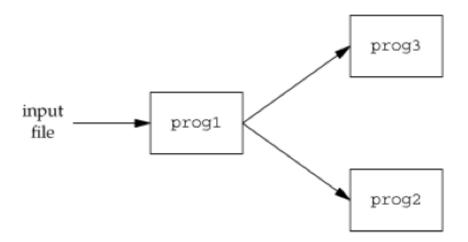
```
#include <sys/stat.h>
int mkfifo(const char *path, mode_t mode);

Returns: 0 if OK, -1 otherwise
```

- aka "named pipes"
- allows unrelated processes to communicate
- just a type of file test for using S\_ISFIFO(st\_mode)
- mode same as for open(2)
- use regular I/O operations (ie open(2), read(2), write(2), unlink(2) etc.)
- used by shell commands to pass data from one shell pipeline to another without creating intermediate temporary files

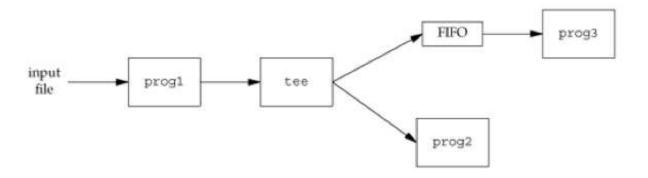
# FIFOs: mkfifo(2)

# Example: split input into sets



# FIFOs: mkfifo(2)

#### Example: split input into sets



- \$ mkfifo fifo
- \$ grep pattern fifo > match &
- \$ gzcat file.gz | tee fifo | grep -v pattern > nomatch

# System V IPC

Three types of IPC originating from System V:

- Semaphores
- Shared Memory
- Message Queues

All three use *IPC structures*, referred to by an *identifier* and a *key*; all three are (necessarily) limited to communication between processes on one and the same host.

Since these structures are not known by name, special system calls (msgget(2), semop(2), shmat(2), etc.) and special userland commands (ipcrm(1), ipcs(1), etc.) are necessary.

# System V IPC: Semaphores

A semaphore is a counter used to provide access to a shared data object for multiple processes. To obtain a shared resource a process needs to do the following:

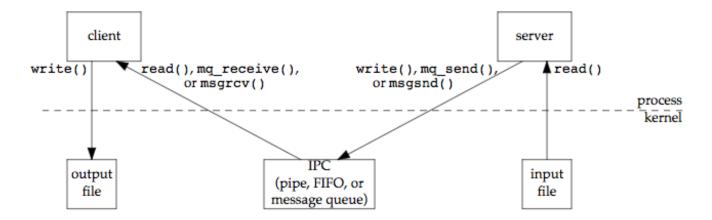
- 1. Test semaphore that controls the resource.
- 2. If value of semaphore > 0, decrement semaphore and use resource; increment semaphore when done
- 3. If value == 0 sleep until value > 0

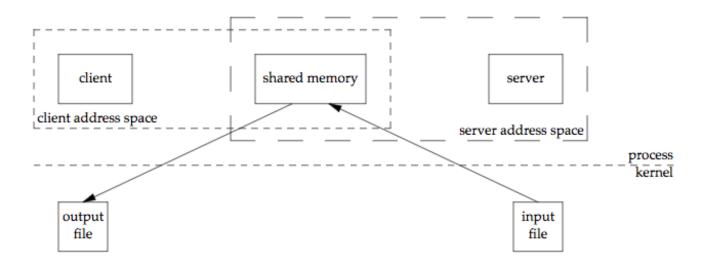
Semaphores are obtained using semget(2), properties controlled using semctl(2), operations on a semaphore performed using semop(2).

# System V IPC: Semaphores

```
$ cc -Wall semdemo.c
1$ ./a.out
2$ ./a.out
$ ipcs -s
```

# IPC data flow

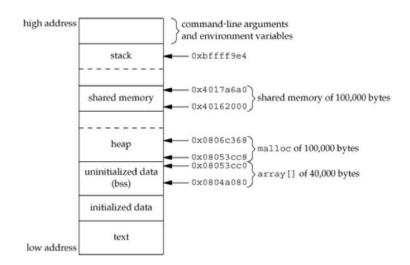




- fastest form of IPC
- access to shared region of memory often controlled using semaphores
- obtain a shared memory identifier using shmget (2)
- catchall for shared memory operations: shmctl(2)
- attach shared memory segment to a processes address space by calling shmat(2)
- detach it using shmdt(2)

```
$ cc -Wall shmdemo.c
$ ./a.out "Cow says: 'Moo!'"
$ ./a.out
$ ipcs -m
```

```
$ cc -Wall memory-layout.c
$ ./a.out
array[] from 804a080 to 8053cc0
stack around bffff9e4
malloced from 8053cc8 to 806c368
shared memory attached from 40162000 to 4017a6a0
```



# System V IPC: Message Queues

- linked list of messages stored in the kernel
- create or open existing queue using msgget (2)
- add message at end of queue using msgsnd(2)
- control queue properties using msgct1(2)
- receive messages from queue using msgrcv(2)

The message itself is contained in a user-defined structure such as

```
struct mymsg {
    long mtype;     /* message type */
    char mtext[512]; /* body of message */
};
```

# System V IPC: Message Queues

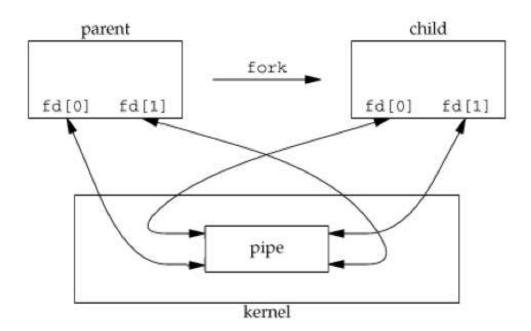
```
$ cc -Wall msgsend.c -o msgsend
$ cc -Wall msgrecv.c -o msgrecv
$ ipcs -q
 ./msgsend 1
$ ipcs -q
$ ./msgsend 1
$ ipcs -q
 ./msgrecv 1
$ ipcs -q
$ ./msgrecv 1
$ ipcs -q
$ ./msgrecv 1
^C
$ ipcs -q
$ ./msgsend 2
$ ipcrm -q <msqid>
```

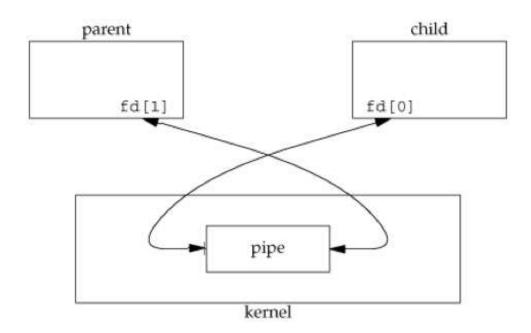
```
#include <sys/socket.h>
int socketpair(int d, int type, int protocol, int *sv);
```

The socketpair(2) call creates an unnamed pair of connected sockets in the specified domain d, of the specified *type*, and using the optionally specified *protocol*.

The descriptors used in referencing the new sockets are returned in sv[0] and sv[1]. The two sockets are indistinguishable.

This call is currently implemented only for the UNIX domain.





```
$ cc -Wall socketpair.c
$ ./a.out
78482 --> sending: In Xanadu, did Kublai Khan . . .
78483 --> sending: A stately pleasure dome decree . . .
78483 --> reading: In Xanadu, did Kublai Khan . . .
78482 --> reading: A stately pleasure dome decree . . .
$
```

# Sockets: socket(2)

```
#include <sys/socket.h>
int socket(int domain, int type, int protocol);
```

#### Some of the currently supported domains are:

Domain	Description
PF_LOCAL	local (previously UNIX) domain protocols
PF_INET	ARPA Internet protocols
PF_INET6	ARPA IPv6 (Internet Protocol version 6) protocols
PF_ARP	RFC 826 Ethernet Address Resolution Protocol

#### Some of the currently defined types are:

Type	Description
SOCK_STREAM	sequenced, reliable, two-way connection based byte streams
SOCK_DGRAM	connectionless, unreliable messages of a fixed (typically small) maximum length
SOCK_RAW	access to internal network protocols and interfaces

# Sockets: Datagrams in the UNIX/LOCAL domain

```
1$ cc -Wall udgramsend.c -o send
1$ cc -Wall udgramread.c -o read
1$ ./read
socket --> socket

2$ ls -l socket
srwxr-xr-x 1 jans users 0 Oct 31 19:17 socket
2$ ./send socket
2$
--> The sea is calm tonight, the tide is full . . .
1$
```

# Sockets: Datagrams in the UNIX/LOCAL domain

- create socket using socket (2)
- attach to a socket using bind(2)
- binding a name in the UNIX domain creates a socket in the file system
- both processes need to agree on the name to use
- these files are only used for rendezvous, not for message delivery once a connection has been established
- sockets must be removed using unlink(2)

# Sockets: Datagrams in the Internet Domain

# Sockets: Datagrams in the Internet Domain

- Unlike UNIX domain names, Internet socket names are not entered into the file system and, therefore, they do not have to be unlinked after the socket has been closed.
- The local machine address for a socket can be any valid network address of the machine, if it has more than one, or it can be the wildcard value INADDR\_ANY.
- "well-known" ports (range 1 1023) only available to super-user
- request any port by calling bind(2) with a port number of 0
- determine used port number (or other information) using getsockname(2)
- convert between network byteorder and host byteorder using htons(3) and ntohs(3) (which may be noops)

# Sockets: Connections using stream sockets

```
1$ cc -Wall streamread.c -o read
1$ cc -Wall streamwrite.c -o write
1$ ./read
Socket has port #65398
2$ ./write localhost 65398
2$ ./write localhost 65398
--> Half a league, half a league . . .
Ending connection
--> Half a league, half a league . . .
Ending connection
2$ nc localhost 65398
moo
2$
```

# Sockets: Connections using stream sockets

- connections are asymmetrical: one process requests a connection,
   the other process accepts the request
- one socket is created for each accepted request
- mark socket as willing to accept connections using listen(2)
- pending connections are then accept (2) ed
- accept (2) will block if no connections are available
- select(2) to check if connection requests are pending

# Sockets: Connections using stream sockets

```
1$ cc -Wall strchkread.c -o read
1$ ./read
Socket has port #65398
Do something else
Do something else
2$ ./write localhost 65398
2$ ./write localhost 65398
-> Half a league, half a league . . .
Ending connection
Do something else
--> Half a league, half a league . . .
Ending connection
^C
1$
```

#### Sockets: Other Useful Functions

I/O on sockets is done on descriptors, just like regular I/O, ie the typical read(2) and write(2) calls will work. In order to specify certain flags, some other functions can be used:

- send(2), sendto(2) and sendmsg(2)
- recv(2), recvfrom(2) and recvmsg(2)

#### To manipulate the options associated with a socket, use setsockopt(2):

Option	Description
SO_DEBUG	enables recording of debugging information
SO_REUSEADDR	enables local address reuse
SO_REUSEPORT	enables duplicate address and port bindings
SO_KEEPALIVE	enables keep connections alive
SO_DONTROUTE	enables routing bypass for outgoing messages
SO_LINGER	linger on close if data present
SO_BROADCAST	enables permission to transmit broadcast messages
SO_OOBINLINE	enables reception of out-of-band data in band
SO_SNDBUF	set buffer size for output
SO_RCVBUF	set buffer size for input
SO_SNDLOWAT	set minimum count for output
SO_RCVLOWAT	set minimum count for input
SO_SNDTIMEO	set timeout value for output
SO_RCVTIMEO	set timeout value for input
SO_TIMESTAMP	enables reception of a timestamp with datagrams
SO_TYPE	get the type of the socket (get only)
SO_ERROR	get and clear error on the socket (get only)

#### More Information

- http://www.cs.stevens.edu/~jschauma/631/ipctut.pdf
- http://www.cs.stevens.edu/~jschauma/631/ipc.pdf
- http://www.cs.cf.ac.uk/Dave/C/node25.html
- http://kohala.com/start/unpv22e/unpv22e.chap12.pdf

http://beej.us/guide/bgipc/output/html/singlepage/bgipc.html

HW#3: write the basic framework for your final project.

https://www.cs.stevens.edu/~jschauma/631/f16-hw3.html