

Answer: Kerckhoffs' Principle is a concept in cryptography that states a cryptographic system should be secure even if everything about the system, except the key, is public knowledge. This principle emphasizes that the security of a system should not rely on its algorithm being kept secret, but rather on the secrecy of the key used to encrypt and decrypt messages.

## Modular arithmetic

- $x = y \bmod N$  if and only if  $N$  divides  $x - y$
- $[x \bmod N]$  = the remainder when  $x$  is divided by  $N$ 
  - I.e., the unique value  $y \in \{0, \dots, N-1\}$  such that  $x = y \bmod N$
- $25 = 35 \bmod 10$
- $25 \neq [35 \bmod 10]$
- $5 = [35 \bmod 10]$

# The Vigenère cipher

- The key is *multiple* characters, not just one
- To encrypt, shift each character in the plaintext by the amount dictated by the next character of the key
  - Wrap around in the key as needed
- Decryption just reverses the process

tellhimaboutme
cafecafecafeca
<b>veq<p>j</p>iredozxoe</b>

## Attacking the Vigenère cipher

- Look at every 14<sup>th</sup> character of the ciphertext, starting with the first
  - Call this the first “stream”
- Let  $\alpha$  be the most common character appearing in this stream
- Most likely,  $\alpha$  corresponds to the most common character of the plaintext (i.e., ‘e’)
  - Guess that the first character of the key is  $\alpha - 'e'$
- Repeat for all other positions

## A better attack (high level)

- Let  $p_i$  ( $0 \leq i \leq 25$ ) denote the frequency of the  $i^{\text{th}}$  English letter in normal English plaintext
  - One can compute that  $\sum_i p_i^2 \approx 0.065$
- Let  $q_i$  denote the observed frequency of the  $i^{\text{th}}$  letter in a given stream of the ciphertext
- If the shift for that stream is  $j$ , expect  $q_{i+j} \approx p_i$  for all  $i$ 
  - So expect  $\sum_i p_i q_{i+j} \approx 0.065$
- Test for every value of  $j$  to find the right one
  - Repeat for each stream

## Threat models for encryption

- Ciphertext-only attack
  - One ciphertext or many?
- Known-plaintext attack
- Chosen-plaintext attack
- Chosen-ciphertext attack

### 1. 唯密文攻击 (Ciphertext-Only Attack)

- **定义：** 攻击者仅能获取加密后的密文，但不知道对应的明文或密钥。这是最基础的攻击场景。

### 2. 已知明文攻击 (Known-Plaintext Attack)

- **定义：** 攻击者掌握部分明文及其对应的密文，目标是破解密钥或解密其他密文。

### 3. 选择明文攻击 (Chosen-Plaintext Attack)

- **定义：** 攻击者可以主动选择任意明文，并获取对应的密文，目标是推断密钥或解密其他密文。

### 4. 选择密文攻击 (Chosen-Ciphertext Attack)

- **定义：** 攻击者可以提交任意密文，并获取解密后的明文，目标是破解密钥或伪造合法密文。

## Core principles of modern crypto

- **Formal definitions**

- Precise, mathematical model and definition of what security means

- **Assumptions**

- Clearly stated and unambiguous

- **Proofs of security**

- Move away from design-break-patch cycle

## The right definition

- “Regardless of any *prior* information the attacker has about the plaintext, the ciphertext should leak no *additional* information about the plaintext”