## Australian National University Research School of Computer Science

## COMP3600/COMP6466 in 2016 - Quiz four

Due: 23:55pm Friday, October 14

Submit your work electronically through Wattle. The total mark of this quiz worths 20 points which is worth 4 of the final mark.

## Question 1 (10 points).

- (a) Apply the forest consisting of (inverted) directed tree implementation of disjoint-set data structure, using both heuristics (path compression and union by rank), to find the connected components in a graph with 10 vertices and edges provided in this order: 1-3, 2-5, 3-8, 4-7, 6-7, 8-9, 1-10. (3 points)
- (b) Consider a red-black tree formed by inserting n nodes with RB-INSERT. Argue if n > 1, the tree has at least one red node. (2 points)
- (c) Given an algorithm that determines whether or not a given undirected graph G = (V, E) contains a cycle. Your algorithm should be run in O(|V|) time, independent of |E|. (2 points)
- (d) Given a connected undirected graph G = (V, E) where V is the set of nodes and E is the set of edges, devise an O(|V| + |E|) algorithm to verify whether G contains any odd cycles. An odd cycle in a graph is a simple cycle that has odd numbers of edges. (3 points).

## Question 2 (10 points).

- (a) In which case should the Bellman-Ford algorithm instead of Dijkstra's algorithm be applied to solve the single-source shortest paths problem? (1 point)
- (b) Given a connected undirected graph G = (V, E), assume that each edge  $e \in E$  has a non-negative weight, let  $e_{min}$  be an edge with the minimum weight in a cycle of G, prove or disprove that  $e_{min}$  will be contained by any minimum spanning tree in G. **Hints:** you may just give a counter-example if you disprove the claim. (3 points)

- (c) You are given a set of cities, along with the pattern of highways between them, in the form of an undirected graph G = (V, E). Each stretch of highway  $e \in E$  connects two of the cities, and you know its length in kilometers l(e). You want to get from city  $s \in V$  to city  $t \in V$ . There is one problem: your car can only hold enough petrol to cover L kilometers. There are petrol stations in each city, but not between cities. Therefore, you can only take a route if every one of its edges has length no greater than L.
  - [i] Given the limitation on your car's fuel tank capacity, show how to determine in linear time whether there is a feasible route from s to t. (3 points)
  - [ii] You are now planning to buy a new car, and you want to know the minimum fuel tank capacity that is needed to travel from s to t. Give an  $O((|V|+|E|)\log |V|)$  algorithm to determine this. (3 points)