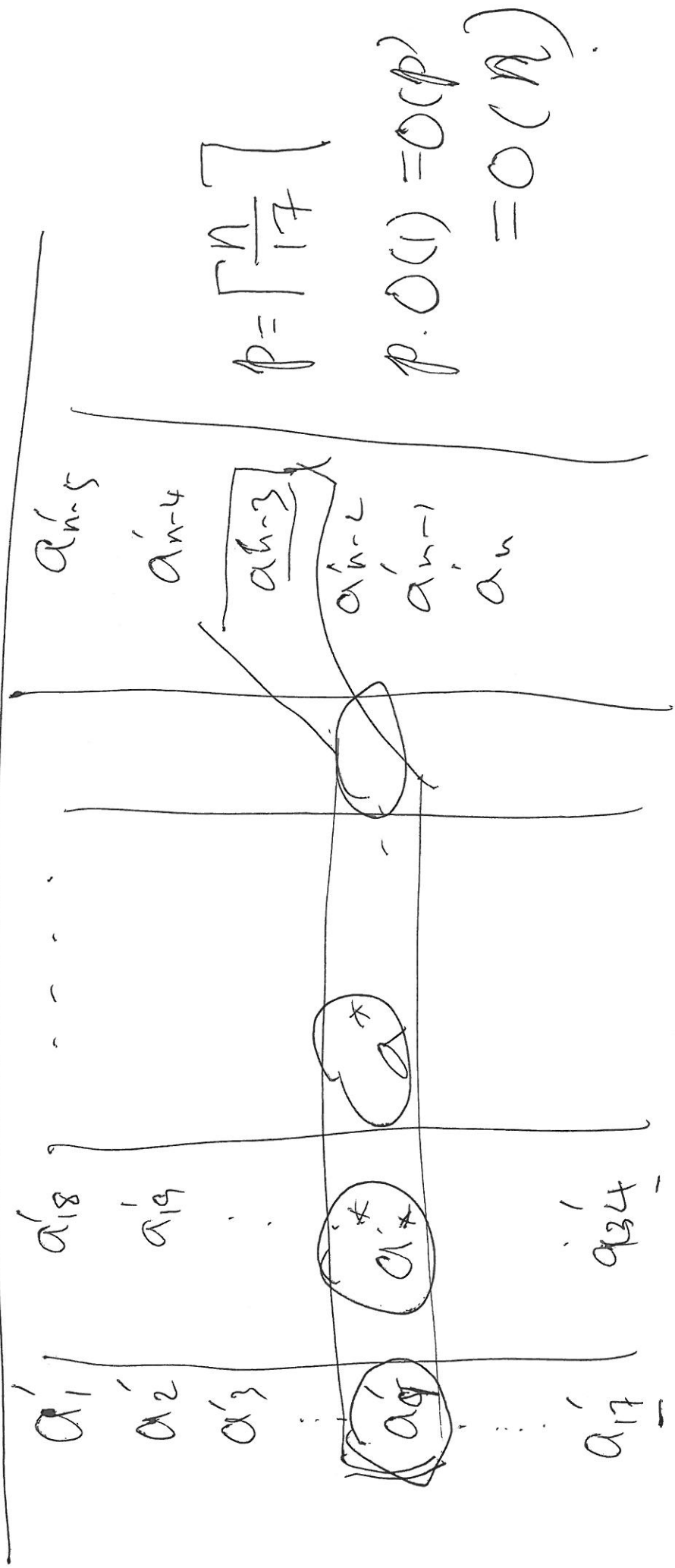


1. Divide & Conquer.

2. Dynamic Programming

3. Greedy Policy.



$$p = \lceil \frac{n}{17} \rceil$$

$$p \cdot O(1) = O(p) = O(n)$$

g_1, g_2, \dots, g_p

• Median Sequence. $\lceil \frac{n}{17} \rceil$ $T(n)$ 8

$a_1^+ \ a_1^+ \ \dots \ a_{n-3}$ the median \mathcal{L}

$$R_1 = \{a_i \mid a_i < x\} \quad T\left(\frac{n}{17}\right)$$

$$R_2 = \{a_i \mid a_i = x\} \quad O(n)$$

$$R_3 = \{a_i \mid a_i > x\}$$

$$\max \{T(R_1), T(R_3)\}$$

$$|R_1| \leq ? \quad |R_3| \leq ?$$

④

$$\frac{|R_3|}{1} \geq 9 \left(\frac{\lceil \frac{n}{17} \rceil}{2} - 2 \right)$$

$$\geq 9 \left(\frac{n}{2 \times 17} - 2 \right)$$

$$\frac{|R_3|}{1} \geq \frac{9n}{34} - 18$$

$$|R_1| = n - |A| - |B| - |R_3| \leq n - |R_3|$$

$$\leq n - \left(\frac{9n}{34} - 18 \right)$$

$$= \frac{(34-9)n}{34} + 18 = \frac{25n}{34} + 18$$

$$T(n) = \overline{T\left(\frac{n}{17}\right) + T\left(\left\lceil \frac{n}{17} \right\rceil\right) + cn} + \max(T(R_1), T(R_2))$$

$$T(n) = T\left(\left\lceil \frac{n}{17} \right\rceil\right) + T\left(\frac{25n}{34} + 18\right) + cn$$

1 spell

Edit distance

1 match

-1 mismatch

-2 for a gap

Snowy
Sunny

S - N O W Y
S U N N - Y

1 -2 1 -1 -2 1 = -2

- S N O W - Y
S U N - - N Y

-2 -1 1 -2 -2 -2 1

= -7

②

$$X_i = x_1 \ x_2 \ \dots \ x_n$$

$$Y_i = y_1 \ y_2 \ \dots \ y_n$$

$$i=0, \text{ } j=0$$

$$C(i, j) =$$

$$\begin{cases} C(i-1, j-1) + 1 & X_i = Y_j \\ \max \{ C(i-1, j), C(i, j-1) \} - 1 & X_i \neq Y_j \end{cases}$$

align with gap? align with



$$O(m \cdot n) \quad \underbrace{C(m, n)}_{X_m \ Y_n}$$

Index matrix

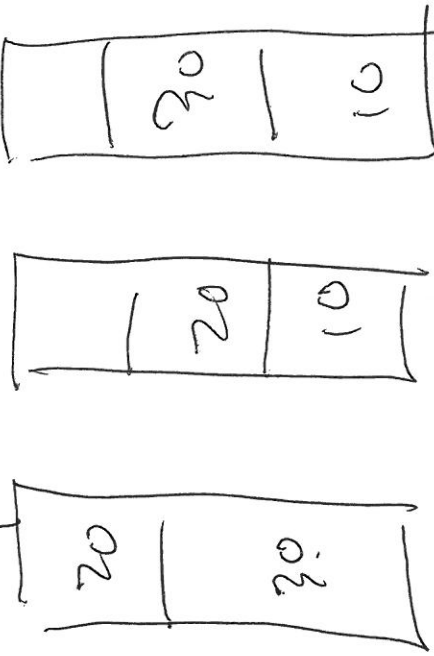
Container 50 kilo

Item Price, weight

P_i W_i

	item 1	item 2	item 3
w_i	10	20 kilo	30 kilo
P_i	\$60	\$100	\$120

~~$\frac{P_i}{W_i}$~~



\$220 \$160 \$180

item 1	item 2	item 3
6	5	4
10	20	20
\$60	\$100	\$80
\$240		

7

(8)

Contradiction method

Assume that this solution is not optimal.

Let $\underline{I}_1, \underline{I}_2, \dots, \underline{I}_k$ be the optimal solution.

$\underline{J}_1, \underline{J}_2, \dots, \underline{J}_k$ — our solution

p_i, w_i . \underline{J}_k in our solution
— be the 1st item is not in the optimal solution

$$\text{OPT}' = \text{OPT} + \{ \underline{J}_k \} - \frac{\underline{I}_k(w_k)}{w_k} - \frac{p_k * w_k}{w_k} \geq 0$$

$p_k * w_k$