

The Standard Libraries

Information technology — Programming Languages — Ada

Section 1: General

Ada is a programming language designed to support the construction of long-lived, highly reliable software systems. The language includes facilities to define packages of related types, objects, and operations. The packages may be parameterized and the types may be extended to support the construction of libraries of reusable, adaptable software components. The operations may be implemented as subprograms using conventional sequential control structures, or as entries that include synchronization of concurrent threads of control as part of their invocation. The language treats modularity in the physical sense as well, with a facility to support separate compilation.

The language includes a complete facility for the support of real-time, concurrent programming. Errors can be signaled as exceptions and handled explicitly. The language also covers systems programming; this requires precise control over the representation of data and access to system-dependent properties. Finally, a predefined environment of standard packages is provided, including facilities for, among others, input-output, string manipulation, numeric elementary functions, and random number generation.

1.1 Scope

This International Standard specifies the form and meaning of programs written in Ada. Its purpose is to promote the portability of Ada programs to a variety of data processing systems.

1.1.1 Extent

This International Standard specifies:

- The form of a program written in Ada;
- The effect of translating and executing such a program;
- The manner in which program units may be combined to form Ada programs;
- The language-defined library units that a conforming implementation is required to supply;
- The permissible variations within the standard, and the manner in which they are to be documented;

- 7 • Those violations of the standard that a conforming implementation is required to detect, and the
effect of attempting to translate or execute a program containing such violations;
- 8 • Those violations of the standard that a conforming implementation is not required to detect.

9 This International Standard does not specify:

- 10 • The means whereby a program written in Ada is transformed into object code executable by a
processor;
- 11 • The means whereby translation or execution of programs is invoked and the executing units are
controlled;
- 12 • The size or speed of the object code, or the relative execution speed of different language
constructs;
- 13 • The form or contents of any listings produced by implementations; in particular, the form or
contents of error or warning messages;
- 14 • The effect of unspecified execution.
- 15 • The size of a program or program unit that will exceed the capacity of a particular conforming
implementation.

1.1.2 Structure

1 This International Standard contains thirteen sections, fourteen annexes, and an index.

2 The *core* of the Ada language consists of:

- 3 • Sections 1 through 13
- 4 • Annex A, “Predefined Language Environment”
- 5 • Annex B, “Interface to Other Languages”
- 6 • Annex J, “Obsolescent Features”

7 The following *Specialized Needs Annexes* define features that are needed by certain application areas:

- 8 • Annex C, “Systems Programming”
- 9 • Annex D, “Real-Time Systems”
- 10 • Annex E, “Distributed Systems”
- 11 • Annex F, “Information Systems”
- 12 • Annex G, “Numerics”
- 13 • Annex H, “High Integrity Systems”

14 The core language and the Specialized Needs Annexes are normative, except that the material in each of
the items listed below is informative:

- 15 • Text under a NOTES or Examples heading.
- 16 • Each clause or subclause whose title starts with the word “Example” or “Examples”.

17 All implementations shall conform to the core language. In addition, an implementation may conform
separately to one or more Specialized Needs Annexes.

The following Annexes are informative:

• Annex K, “Language-Defined Attributes”	18
• Annex L, “Language-Defined Pragmas”	19
• M.2, “Implementation-Defined Characteristics”	20
• Annex N, “Glossary”	21
• Annex P, “Syntax Summary”	22
	23

Each section is divided into clauses and subclauses that have a common structure. Each section, clause, and subclause first introduces its subject. After the introductory text, text is labeled with the following headings:

<i>Syntax</i>	
Syntax rules (indented).	25
<i>Name Resolution Rules</i>	
Compile-time rules that are used in name resolution, including overload resolution.	26
<i>Legality Rules</i>	
Rules that are enforced at compile time. A construct is <i>legal</i> if it obeys all of the Legality Rules.	27
<i>Static Semantics</i>	
A definition of the compile-time effect of each construct.	28
<i>Post-Compilation Rules</i>	
Rules that are enforced before running a partition. A partition is legal if its compilation units are legal and it obeys all of the Post-Compilation Rules.	29
<i>Dynamic Semantics</i>	
A definition of the run-time effect of each construct.	30
<i>Bounded (Run-Time) Errors</i>	
Situations that result in bounded (run-time) errors (see 1.1.5).	31
<i>Erroneous Execution</i>	
Situations that result in erroneous execution (see 1.1.5).	32
<i>Implementation Requirements</i>	
Additional requirements for conforming implementations.	33
<i>Documentation Requirements</i>	
Documentation requirements for conforming implementations.	34
<i>Metrics</i>	
Metrics that are specified for the time/space properties of the execution of certain language constructs.	35
<i>Implementation Permissions</i>	
Additional permissions given to the implementer.	36

Implementation Advice

- 37 Optional advice given to the implementer. The word “should” is used to indicate that the advice is a recommendation, not a requirement. It is implementation defined whether or not a given recommendation is obeyed.

NOTES

- 38 1 Notes emphasize consequences of the rules described in the (sub)clause or elsewhere. This material is informative.

Examples

- 39 Examples illustrate the possible forms of the constructs described. This material is informative.

1.1.3 Conformity of an Implementation with the Standard

Implementation Requirements

- 1 A conforming implementation shall:
- 2 • Translate and correctly execute legal programs written in Ada, provided that they are not so large as to exceed the capacity of the implementation;
 - 3 • Identify all programs or program units that are so large as to exceed the capacity of the implementation (or raise an appropriate exception at run time);
 - 4 • Identify all programs or program units that contain errors whose detection is required by this International Standard;
 - 5 • Supply all language-defined library units required by this International Standard;
 - 6 • Contain no variations except those explicitly permitted by this International Standard, or those that are impossible or impractical to avoid given the implementation's execution environment;
 - 7 • Specify all such variations in the manner prescribed by this International Standard.
- 8 The *external effect* of the execution of an Ada program is defined in terms of its interactions with its external environment. The following are defined as *external interactions*:
- 9 • Any interaction with an external file (see A.7);
 - 10 • The execution of certain `code_statements` (see 13.8); which `code_statements` cause external interactions is implementation defined.
 - 11 • Any call on an imported subprogram (see Annex B), including any parameters passed to it;
 - 12 • Any result returned or exception propagated from a main subprogram (see 10.2) or an exported subprogram (see Annex B) to an external caller;
 - 13 • Any read or update of an atomic or volatile object (see C.6);
 - 14 • The values of imported and exported objects (see Annex B) at the time of any other interaction with the external environment.
- 15 A conforming implementation of this International Standard shall produce for the execution of a given Ada program a set of interactions with the external environment whose order and timing are consistent with the definitions and requirements of this International Standard for the semantics of the given program.
- 16 An implementation that conforms to this Standard shall support each capability required by the core language as specified. In addition, an implementation that conforms to this Standard may conform to one or more Specialized Needs Annexes (or to none). Conformance to a Specialized Needs Annex means that each capability required by the Annex is provided as specified.

An implementation conforming to this International Standard may provide additional attributes, library units, and pragmas. However, it shall not provide any attribute, library unit, or pragma having the same name as an attribute, library unit, or pragma (respectively) specified in a Specialized Needs Annex unless the provided construct is either as specified in the Specialized Needs Annex or is more limited in capability than that required by the Annex. A program that attempts to use an unsupported capability of an Annex shall either be identified by the implementation before run time or shall raise an exception at run time.

Documentation Requirements

Certain aspects of the semantics are defined to be either *implementation defined* or *unspecified*. In such cases, the set of possible effects is specified, and the implementation may choose any effect in the set. Implementations shall document their behavior in implementation-defined situations, but documentation is not required for unspecified situations. The implementation-defined characteristics are summarized in M.2.

The implementation may choose to document implementation-defined behavior either by documenting what happens in general, or by providing some mechanism for the user to determine what happens in a particular case.

Implementation Advice

If an implementation detects the use of an unsupported Specialized Needs Annex feature at run time, it should raise Program_Error if feasible.

If an implementation wishes to provide implementation-defined extensions to the functionality of a language-defined library unit, it should normally do so by adding children to the library unit.

NOTES

- 2 The above requirements imply that an implementation conforming to this Standard may support some of the capabilities required by a Specialized Needs Annex without supporting all required capabilities.

1.1.4 Method of Description and Syntax Notation

The form of an Ada program is described by means of a context-free syntax together with context-dependent requirements expressed by narrative rules.

The meaning of Ada programs is described by means of narrative rules defining both the effects of each construct and the composition rules for constructs.

The context-free syntax of the language is described using a simple variant of Backus-Naur Form. In particular:

- Lower case words in a sans-serif font, some containing embedded underlines, are used to denote syntactic categories, for example:

`case_statement`

- Boldface words are used to denote reserved words, for example:

`array`

- Square brackets enclose optional items. Thus the two following rules are equivalent.

`simple_return_statement ::= return [expression];`

`simple_return_statement ::= return; | return expression;`

- 10 • Curly brackets enclose a repeated item. The item may appear zero or more times; the repetitions occur from left to right as with an equivalent left-recursive rule. Thus the two following rules are equivalent.

11 term ::= factor {multiplying_operator factor}
 term ::= factor | term multiplying_operator factor

- 12 • A vertical line separates alternative items unless it occurs immediately after an opening curly bracket, in which case it stands for itself:

13 constraint ::= scalar_constraint | composite_constraint
 discrete_choice_list ::= discrete_choice {| discrete_choice}

- 14 • If the name of any syntactic category starts with an italicized part, it is equivalent to the category name without the italicized part. The italicized part is intended to convey some semantic information. For example *subtype_name* and *task_name* are both equivalent to *name* alone.

14.1/2 The delimiters, compound delimiters, reserved words, and *numeric_literals* are exclusively made of the characters whose code position is between 16#20# and 16#7E#, inclusively. The special characters for which names are defined in this International Standard (see 2.1) belong to the same range. For example, the character E in the definition of exponent is the character whose name is “LATIN CAPITAL LETTER E”, not “GREEK CAPITAL LETTER EPSILON”.

14.2/2 When this International Standard mentions the conversion of some character or sequence of characters to upper case, it means the character or sequence of characters obtained by using locale-independent full case folding, as defined by documents referenced in the note in section 1 of ISO/IEC 10646:2003.

15 A *syntactic category* is a nonterminal in the grammar defined in BNF under “Syntax.” Names of syntactic categories are set in a different font, *like_this*.

16 A *construct* is a piece of text (explicit or implicit) that is an instance of a syntactic category defined under “Syntax”.

17 A *constituent* of a construct is the construct itself, or any construct appearing within it.

18 Whenever the run-time semantics defines certain actions to happen in an *arbitrary order*, this means that the implementation shall arrange for these actions to occur in a way that is equivalent to some sequential order, following the rules that result from that sequential order. When evaluations are defined to happen in an arbitrary order, with conversion of the results to some subtypes, or with some run-time checks, the evaluations, conversions, and checks may be arbitrarily interspersed, so long as each expression is evaluated before converting or checking its value. Note that the effect of a program can depend on the order chosen by the implementation. This can happen, for example, if two actual parameters of a given call have side effects.

NOTES

19 3 The syntax rules describing structured constructs are presented in a form that corresponds to the recommended paragraphing. For example, an *if_statement* is defined as:

20 if_statement ::=
 if condition then
 sequence_of_statements
 {elsif condition then
 sequence_of_statements}
 [else
 sequence_of_statements]
 end if;

4 The line breaks and indentation in the syntax rules indicate the recommended line breaks and indentation in the corresponding constructs. The preferred places for other line breaks are after semicolons. 21

1.1.5 Classification of Errors

Implementation Requirements

The language definition classifies errors into several different categories:

- Errors that are required to be detected prior to run time by every Ada implementation; 2

These errors correspond to any violation of a rule given in this International Standard, other than those listed below. In particular, violation of any rule that uses the terms shall, allowed, permitted, legal, or illegal belongs to this category. Any program that contains such an error is not a legal Ada program; on the other hand, the fact that a program is legal does not mean, *per se*, that the program is free from other forms of error. 3

The rules are further classified as either compile time rules, or post compilation rules, depending on whether a violation has to be detected at the time a compilation unit is submitted to the compiler, or may be postponed until the time a compilation unit is incorporated into a partition of a program. 4

- Errors that are required to be detected at run time by the execution of an Ada program; 5

The corresponding error situations are associated with the names of the predefined exceptions. Every Ada compiler is required to generate code that raises the corresponding exception if such an error situation arises during program execution. If such an error situation is certain to arise in every execution of a construct, then an implementation is allowed (although not required) to report this fact at compilation time. 6

- Bounded errors; 7

The language rules define certain kinds of errors that need not be detected either prior to or during run time, but if not detected, the range of possible effects shall be bounded. The errors of this category are called *bounded errors*. The possible effects of a given bounded error are specified for each such error, but in any case one possible effect of a bounded error is the raising of the exception `Program_Error`. 8

- Erroneous execution. 9

In addition to bounded errors, the language rules define certain kinds of errors as leading to *erroneous execution*. Like bounded errors, the implementation need not detect such errors either prior to or during run time. Unlike bounded errors, there is no language-specified bound on the possible effect of erroneous execution; the effect is in general not predictable. 10

Implementation Permissions

An implementation may provide *nonstandard modes* of operation. Typically these modes would be selected by a `pragma` or by a command line switch when the compiler is invoked. When operating in a nonstandard mode, the implementation may reject `compilation_units` that do not conform to additional requirements associated with the mode, such as an excessive number of warnings or violation of coding style guidelines. Similarly, in a nonstandard mode, the implementation may apply special optimizations or alternative algorithms that are only meaningful for programs that satisfy certain criteria specified by the implementation. In any case, an implementation shall support a *standard mode* that conforms to the requirements of this International Standard; in particular, in the standard mode, all legal `compilation_units` shall be accepted. 11

Implementation Advice

- 12 If an implementation detects a bounded error or erroneous execution, it should raise Program_Error.

1.2 Normative References

- 1 The following standards contain provisions which, through reference in this text, constitute provisions of this International Standard. At the time of publication, the editions indicated were valid. All standards are subject to revision, and parties to agreements based on this International Standard are encouraged to investigate the possibility of applying the most recent editions of the standards indicated below. Members of IEC and ISO maintain registers of currently valid International Standards.
- 2 ISO/IEC 646:1991, *Information technology — ISO 7-bit coded character set for information interchange*.
- 3/2 ISO/IEC 1539-1:2004, *Information technology — Programming languages — Fortran — Part 1: Base language*.
- 4/2 ISO/IEC 1989:2002, *Information technology — Programming languages — COBOL*.
- 5 ISO/IEC 6429:1992, *Information technology — Control functions for coded graphic character sets*.
- 5.1/2 ISO 8601:2004, *Data elements and interchange formats — Information interchange — Representation of dates and times*.
- 6 ISO/IEC 8859-1:1987, *Information processing — 8-bit single-byte coded character sets — Part 1: Latin alphabet No. 1*.
- 7/2 ISO/IEC 9899:1999, *Programming languages — C*, supplemented by Technical Corrigendum 1:2001 and Technical Corrigendum 2:2004.
- 8/2 ISO/IEC 10646:2003, *Information technology — Universal Multiple-Octet Coded Character Set (UCS)*.
- 9/2 ISO/IEC 14882:2003, *Programming languages — C++*.
- 10/2 ISO/IEC TR 19769:2004, *Information technology — Programming languages, their environments and system software interfaces — Extensions for the programming language C to support new character data types*.

1.3 Definitions

- 1/2 Terms are defined throughout this International Standard, indicated by *italic* type. Terms explicitly defined in this International Standard are not to be presumed to refer implicitly to similar terms defined elsewhere. Mathematical terms not defined in this International Standard are to be interpreted according to the *CRC Concise Encyclopedia of Mathematics, Second Edition*. Other terms not defined in this International Standard are to be interpreted according to the *Webster's Third New International Dictionary of the English Language*. Informal descriptions of some terms are also given in Annex N, “Glossary”.

Section 2: Lexical Elements

The text of a program consists of the texts of one or more compilations. The text of a compilation is a sequence of lexical elements, each composed of characters; the rules of composition are given in this section. Pragmas, which provide certain information for the compiler, are also described in this section.

2.1 Character Set

The character repertoire for the text of an Ada program consists of the entire coding space described by the ISO/IEC 10646:2003 Universal Multiple-Octet Coded Character Set. This coding space is organized in planes, each plane comprising 65536 characters.

Syntax

Paragraphs 2 and 3 were deleted.

A character is defined by this International Standard for each cell in the coding space described by ISO/IEC 10646:2003, regardless of whether or not ISO/IEC 10646:2003 allocates a character to that cell.

Static Semantics

The coded representation for characters is implementation defined (it need not be a representation defined within ISO/IEC 10646:2003). A character whose relative code position in its plane is 16#FFFE# or 16#FFFF# is not allowed anywhere in the text of a program.

The semantics of an Ada program whose text is not in Normalization Form KC (as defined by section 24 of ISO/IEC 10646:2003) is implementation defined.

The description of the language definition in this International Standard uses the character properties General Category, Simple Uppercase Mapping, Uppercase Mapping, and Special Case Condition of the documents referenced by the note in section 1 of ISO/IEC 10646:2003. The actual set of graphic symbols used by an implementation for the visual representation of the text of an Ada program is not specified.

Characters are categorized as follows:

This paragraph was deleted.

letter_uppercase

Any character whose General Category is defined to be “Letter, Uppercase”.

letter_lowercase

Any character whose General Category is defined to be “Letter, Lowercase”.

letter_titlecase

Any character whose General Category is defined to be “Letter, Titlecase”.

letter_modifier

Any character whose General Category is defined to be “Letter, Modifier”.

letter_other

Any character whose General Category is defined to be “Letter, Other”.

mark_non_spacing

Any character whose General Category is defined to be “Mark, Non-Spacing”.

mark_spacing_combining

Any character whose General Category is defined to be “Mark, Spacing Combining”.

- 10/2 **number_decimal**
Any character whose General Category is defined to be “Number, Decimal”.
- 10.1/2 **number_letter**
Any character whose General Category is defined to be “Number, Letter”.
- 10.2/2 **punctuation_connector**
Any character whose General Category is defined to be “Punctuation, Connector”.
- 10.3/2 **other_format**
Any character whose General Category is defined to be “Other, Format”.
- 11/2 **separator_space**
Any character whose General Category is defined to be “Separator, Space”.
- 12/2 **separator_line**
Any character whose General Category is defined to be “Separator, Line”.
- 12.1/2 **separator_paragraph**
Any character whose General Category is defined to be “Separator, Paragraph”.
- 13/2 **format_effector**
The characters whose code positions are 16#09# (CHARACTER TABULATION), 16#0A# (LINE FEED), 16#0B# (LINE TABULATION), 16#0C# (FORM FEED), 16#0D# (CARRIAGE RETURN), 16#85# (NEXT LINE), and the characters in categories **separator_line** and **separator_paragraph**.
- 13.1/2 **other_control**
Any character whose General Category is defined to be “Other, Control”, and which is not defined to be a **format_effector**.
- 13.2/2 **other_private_use**
Any character whose General Category is defined to be “Other, Private Use”.
- 13.3/2 **other_surrogate**
Any character whose General Category is defined to be “Other, Surrogate”.
- 14/2 **graphic_character**
Any character that is not in the categories **other_control**, **other_private_use**, **other_surrogate**, **format_effector**, and whose relative code position in its plane is neither 16#FFE# nor 16#FFF#.
- 15/2 The following names are used when referring to certain characters (the first name is that given in ISO/IEC 10646:2003):

graphic symbol	name	graphic symbol	name
"	quotation mark	:	colon
#	number sign	;	semicolon
&	ampersand	<	less-than sign
'	apostrophe, tick	=	equals sign
(left parenthesis	>	greater-than sign
)	right parenthesis	-	low line, underline
*	asterisk, multiply		vertical line
+	plus sign	/	solidus, divide
,	comma	!	exclamation point
-	hyphen-minus, minus	%	percent sign
.	full stop, dot, point		

*Implementation Permissions**This paragraph was deleted.*

16/2

NOTES

1 The characters in categories other_control, other_private_use, and other_surrogate are only allowed in comments.

17/2

2 The language does not specify the source representation of programs.

18

2.2 Lexical Elements, Separators, and Delimiters*Static Semantics*

The text of a program consists of the texts of one or more compilations. The text of each compilation is a sequence of separate *lexical elements*. Each lexical element is formed from a sequence of characters, and is either a delimiter, an identifier, a reserved word, a numeric_literal, a character_literal, a string_literal, or a comment. The meaning of a program depends only on the particular sequences of lexical elements that form its compilations, excluding comments.

1

The text of a compilation is divided into *lines*. In general, the representation for an end of line is implementation defined. However, a sequence of one or more format_effectors other than the character whose code position is 16#09# (CHARACTER TABULATION) signifies at least one end of line.

2/2

In some cases an explicit *separator* is required to separate adjacent lexical elements. A separator is any of a separator_space, a format_effector, or the end of a line, as follows:

3/2

- A separator_space is a separator except within a comment, a string_literal, or a character_literal.
- The character whose code position is 16#09# (CHARACTER TABULATION) is a separator except within a comment.
- The end of a line is always a separator.

6

One or more separators are allowed between any two adjacent lexical elements, before the first of each compilation, or after the last. At least one separator is required between an identifier, a reserved word, or a numeric_literal and an adjacent identifier, reserved word, or numeric_literal.

7

A *delimiter* is either one of the following characters:

8/2

& ' () * + , - . / : ; < = > |

9

or one of the following *compound delimiters* each composed of two adjacent special characters

10

=> .. ** := /= >= <= << >> <>

11

Each of the special characters listed for single character delimiters is a single delimiter except if this character is used as a character of a compound delimiter, or as a character of a comment, string_literal, character_literal, or numeric_literal.

12

- 13 The following names are used when referring to compound delimiters:

delimiter	name
=>	arrow
..	double dot
**	double star, exponentiate
:=	assignment (pronounced: “becomes”)
/=	inequality (pronounced: “not equal”)
>=	greater than or equal
<=	less than or equal
<<	left label bracket
>>	right label bracket
<>	box

Implementation Requirements

- 14 An implementation shall support lines of at least 200 characters in length, not counting any characters used to signify the end of a line. An implementation shall support lexical elements of at least 200 characters in length. The maximum supported line length and lexical element length are implementation defined.

2.3 Identifiers

- 1 Identifiers are used as names.

Syntax

2/2	identifier ::= identifier_start {identifier_start identifier_extend}
3/2	identifier_start ::= letter_uppercase letter_lowercase letter_titlecase letter_modifier letter_other number_letter
3.1/2	identifier_extend ::= mark_non_spacing mark_spacing_combining number_decimal punctuation_connector other_format

- 4/2 After eliminating the characters in category other_format, an identifier shall not contain two consecutive characters in category punctuation_connector, or end with a character in that category.

Static Semantics

Two identifiers are considered the same if they consist of the same sequence of characters after applying the following transformations (in this order) 5/2

- The characters in category `other_format` are eliminated. 5.1/2
- The remaining sequence of characters is converted to upper case. 5.2/2

After applying these transformations, an identifier shall not be identical to a reserved word (in upper case). 5.3/2

Implementation Permissions

In a nonstandard mode, an implementation may support other upper/lower case equivalence rules for identifiers, to accommodate local conventions. 6

NOTES

3 Identifiers differing only in the use of corresponding upper and lower case letters are considered the same. 6.1/2

Examples

Examples of identifiers:

Count	X	Get_Symbol	Ethelyn	Marion
Snobol_4	X1	Page_Count	Store_Next_Item	
Πλάτων	-- Plato			
Чайковский	-- Tchaikovsky			
θ φ	-- Angles			

2.4 Numeric Literals

There are two kinds of `numeric_literal`, *real literals* and *integer literals*. A real literal is a `numeric_literal` that includes a point; an integer literal is a `numeric_literal` without a point. 1

Syntax

`numeric_literal ::= decimal_literal | based_literal` 2

NOTES

4 The type of an integer literal is *universal_integer*. The type of a real literal is *universal_real*. 3

2.4.1 Decimal Literals

A `decimal_literal` is a `numeric_literal` in the conventional decimal notation (that is, the base is ten). 1

Syntax

`decimal_literal ::= numeral [.numeral] [exponent]` 2

`numeral ::= digit {[underline] digit}` 3

`exponent ::= E [+/-] numeral | E – numeral` 4

`digit ::= 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9` 4.1/2

An exponent for an integer literal shall not have a minus sign. 5

Static Semantics

An underline character in a `numeric_literal` does not affect its meaning. The letter E of an `exponent` can be written either in lower case or in upper case, with the same meaning. 6

An `exponent` indicates the power of ten by which the value of the `decimal_literal` without the `exponent` is to be multiplied to obtain the value of the `decimal_literal` with the `exponent`. 7

Examples

8 Examples of decimal literals:
 9 12 0 1E6 123_456 -- integer literals
 10 12.0 0.0 0.456 3.14159_26 -- real literals

2.4.2 Based Literals

1 A based_literal is a numeric_literal expressed in a form that specifies the base explicitly.

Syntax

2 based_literal ::=
 base # based_numeral [.based_numeral] # [exponent]
 3 base ::= numeral
 4 based_numeral ::=
 extended_digit {[underline] extended_digit}
 5 extended_digit ::= digit | A | B | C | D | E | F

Legality Rules

6 The **base** (the numeric value of the decimal numeral preceding the first #) shall be at least two and at most sixteen. The extended_digits A through F represent the digits ten through fifteen, respectively. The value of each extended_digit of a based_literal shall be less than the base.

Static Semantics

7 The conventional meaning of based notation is assumed. An **exponent** indicates the power of the base by which the value of the based_literal without the exponent is to be multiplied to obtain the value of the based_literal with the exponent. The **base** and the **exponent**, if any, are in decimal notation.
 8 The extended_digits A through F can be written either in lower case or in upper case, with the same meaning.

Examples

9 Examples of based literals:
 10 2#1111_1111# 16#FF# 016#0ff# -- integer literals of value 255
 16#E#E1 2#1110_0000# -- integer literals of value 224
 16#F.FF#E+2 2#1.1111_1111_1110#E11 -- real literals of value 4095.0

2.5 Character Literals

1 A character_literal is formed by enclosing a graphic character between two apostrophe characters.

Syntax

2 character_literal ::= 'graphic_character'
 3 NOTES
 5 A character_literal is an enumeration literal of a character type. See 3.5.2.

*Examples**Examples of character literals:*

'A'	'*'	' '	''	
'L'	'π'	'Λ'		-- Various els.
'∞'	'ꝝ'			-- Big numbers - infinity and aleph.

4

5/2

2.6 String Literals

A `string_literal` is formed by a sequence of graphic characters (possibly none) enclosed between two quotation marks used as string brackets. They are used to represent `operator_symbols` (see 6.1), values of a string type (see 4.2), and array subaggregates (see 4.3.3).

Syntax

```
string_literal ::= "{string_element}"
string_element ::= "" | non_quotation_mark_graphic_character
```

A `string_element` is either a pair of quotation marks (""), or a single `graphic_character` other than a quotation mark.

Static Semantics

The *sequence of characters* of a `string_literal` is formed from the sequence of `string_elements` between the bracketing quotation marks, in the given order, with a `string_element` that is "" becoming a single quotation mark in the sequence of characters, and any other `string_element` being reproduced in the sequence.

A *null string literal* is a `string_literal` with no `string_elements` between the quotation marks.

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NOTES

6 An end of line cannot appear in a `string_literal`.

7 No transformation is performed on the sequence of characters of a `string_literal`.

7

7.1/2

*Examples**Examples of string literals:*

```
"Message of the day:"  
"  
" " "A" " " " -- a null string literal  
-- three string literals of length 1  
  
"Characters such as $, %, and } are allowed in string literals"  
"Archimedes said ""Εύρηκα!""  
"Volume of cylinder ( $\pi r^2 h$ ) = "
```

8

9/2

2.7 Comments

A comment starts with two adjacent hyphens and extends up to the end of the line.

1

Syntax

```
comment ::= --{non_end_of_line_character}  
A comment may appear on any line of a program.
```

2

3

Static Semantics

- 4 The presence or absence of comments has no influence on whether a program is legal or illegal. Furthermore, comments do not influence the meaning of a program; their sole purpose is the enlightenment of the human reader.

Examples

- 5 Examples of comments:

6 -- the last sentence above echoes the Algol 68 report

end; -- processing of Line is complete

-- a long comment may be split onto

-- two or more consecutive lines

----- the first two hyphens start the comment

2.8 Pragmas

- 1 A pragma is a compiler directive. There are language-defined pragmas that give instructions for optimization, listing control, etc. An implementation may support additional (implementation-defined) pragmas.

Syntax

2 `pragma ::= pragma identifier [(pragma_argument_association {, pragma_argument_association})];`

3 `pragma_argument_association ::= [pragma_argument_identifier =>] name
| [pragma_argument_identifier =>] expression`

- 4 In a `pragma`, any `pragma_argument_associations` without a `pragma_argument_identifier` shall precede any associations with a `pragma_argument_identifier`.

- 5 Pragmas are only allowed at the following places in a program:

- 6 • After a semicolon delimiter, but not within a `formal_part` or `discriminant_part`.
- 7 • At any place where the syntax rules allow a construct defined by a syntactic category whose name ends with "declaration", "statement", "clause", or "alternative", or one of the syntactic categories `variant` or `exception_handler`; but not in place of such a construct.
Also at any place where a `compilation_unit` would be allowed.

- 8 Additional syntax rules and placement restrictions exist for specific pragmas.

- 9 The `name` of a `pragma` is the identifier following the reserved word `pragma`. The `name` or `expression` of a `pragma_argument_association` is a `pragma argument`.

- 10 An `identifier specific to a pragma` is an identifier that is used in a `pragma argument` with special meaning for that `pragma`.

Static Semantics

- 11 If an implementation does not recognize the name of a `pragma`, then it has no effect on the semantics of the program. Inside such a `pragma`, the only rules that apply are the Syntax Rules.

Dynamic Semantics

Any **pragma** that appears at the place of an executable construct is executed. Unless otherwise specified for a particular **pragma**, this execution consists of the evaluation of each evaluable **pragma** argument in an arbitrary order.

Implementation Requirements

The implementation shall give a warning message for an unrecognized **pragma** name.

12

Implementation Permissions

An implementation may provide implementation-defined pragmas; the name of an implementation-defined **pragma** shall differ from those of the language-defined pragmas.

14

An implementation may ignore an unrecognized **pragma** even if it violates some of the Syntax Rules, if detecting the syntax error is too complex.

15

Implementation Advice

Normally, implementation-defined pragmas should have no semantic effect for error-free programs; that is, if the implementation-defined pragmas are removed from a working program, the program should still be legal, and should still have the same semantics.

16

Normally, an implementation should not define pragmas that can make an illegal program legal, except as follows:

17

- A **pragma** used to complete a declaration, such as a **pragma Import**;
- A **pragma** used to configure the environment by adding, removing, or replacing **library_items**.

18

19

Syntax

The forms of List, Page, and Optimize **pragmas** are as follows:

20

pragma List(identifier);

21

pragma Page;

22

pragma Optimize(identifier);

23

Other pragmas are defined throughout this International Standard, and are summarized in Annex L.

24

Static Semantics

A **pragma** List takes one of the identifiers On or Off as the single argument. This **pragma** is allowed anywhere a **pragma** is allowed. It specifies that listing of the compilation is to be continued or suspended until a List **pragma** with the opposite argument is given within the same compilation. The **pragma** itself is always listed if the compiler is producing a listing.

25

A **pragma** Page is allowed anywhere a **pragma** is allowed. It specifies that the program text which follows the **pragma** should start on a new page (if the compiler is currently producing a listing).

26

A **pragma** Optimize takes one of the identifiers Time, Space, or Off as the single argument. This **pragma** is allowed anywhere a **pragma** is allowed, and it applies until the end of the immediately enclosing declarative region, or for a **pragma** at the place of a **compilation_unit**, to the end of the **compilation**. It gives advice to the implementation as to whether time or space is the primary optimization criterion, or that optional optimizations should be turned off. It is implementation defined how this advice is followed.

27

Examples

28 Examples of pragmas:

```
pragma List(Off); -- turn off listing generation
pragma Optimize(Off); -- turn off optional optimizations
pragma Inline(Set_Mask); -- generate code for Set_Mask inline
pragma Import(C, Put_Char, External_Name => "putchar"); -- import C putchar function
```

2.9 Reserved Words

Syntax

1/1 This paragraph was deleted.

2/2 The following are the *reserved words*. Within a program, some or all of the letters of a reserved word may be in upper case, and one or more characters in category other_format may be inserted within or at the end of the reserved word.

abort	else	new	return
abs	elsif	not	reverse
abstract	end	null	select
accept	entry	of	separate
access	exception	or	subtype
aliased	exit	others	synchronized
all	for	out	tagged
and	function	overriding	task
array	generic	package	terminate
at	goto	pragma	then
begin	if	private	type
body	in	procedure	until
case	interface	protected	use
constant	is	raise	when
declare	limited	range	while
delay	loop	record	with
delta	mod	rem	
digits		renames	xor
do		requeue	

NOTES

3 8 The reserved words appear in **lower case boldface** in this International Standard, except when used in the designator of an attribute (see 4.1.4). Lower case boldface is also used for a reserved word in a string_literal used as an operator_symbol. This is merely a convention — programs may be written in whatever typeface is desired and available.

Section 3: Declarations and Types

This section describes the types in the language and the rules for declaring constants, variables, and named numbers.

3.1 Declarations

The language defines several kinds of named *entities* that are declared by declarations. The entity's *name* is defined by the declaration, usually by a `defining_identifier`, but sometimes by a `defining_character_literal` or `defining_operator_symbol`.

There are several forms of declaration. A `basic_declaration` is a form of declaration defined as follows.

<i>Syntax</i>	
<code>basic_declaration ::=</code>	
<code>type_declaration</code>	<code> subtype_declaration</code>
<code> object_declaration</code>	<code> number_declaration</code>
<code> subprogram_declaration</code>	<code> abstract_subprogram_declaration</code>
<code> null_procedure_declaration</code>	<code> package_declaration</code>
<code> renaming_declaration</code>	<code> exception_declaration</code>
<code> generic_declaration</code>	<code> generic_instantiation</code>
<code>defining_identifier ::= identifier</code>	

A *declaration* is a language construct that associates a name with (a view of) an entity. A declaration may appear explicitly in the program text (an *explicit declaration*), or may be supposed to occur at a given place in the text as a consequence of the semantics of another construct (an *implicit declaration*).

Each of the following is defined to be a declaration: any `basic_declaration`; an `enumeration_literal_specification`; a `discriminant_specification`; a `component_declaration`; a `loop_parameter_specification`; a `parameter_specification`; a `subprogram_body`; an `entry_declaration`; an `entry_index_specification`; a `choice_parameter_specification`; a `generic_formal_parameter_declaration`. In addition, an `extended_return_statement` is a declaration of its `defining_identifier`.

All declarations contain a *definition* for a *view* of an entity. A view consists of an identification of the entity (the entity *of* the view), plus view-specific characteristics that affect the use of the entity through that view (such as mode of access to an object, formal parameter names and defaults for a subprogram, or visibility to components of a type). In most cases, a declaration also contains the definition for the entity itself (a `renaming_declaration` is an example of a declaration that does not define a new entity, but instead defines a view of an existing entity (see 8.5)).

For each declaration, the language rules define a certain region of text called the *scope* of the declaration (see 8.2). Most declarations associate an identifier with a declared entity. Within its scope, and only there, there are places where it is possible to use the identifier to refer to the declaration, the view it defines, and the associated entity; these places are defined by the visibility rules (see 8.3). At such places the identifier is said to be a *name* of the entity (the `direct_name` or `selector_name`); the name is said to *denote* the declaration, the view, and the associated entity (see 8.6). The declaration is said to *declare* the name, the view, and in most cases, the entity itself.

- 9 As an alternative to an identifier, an enumeration literal can be declared with a character_literal as its name (see 3.5.1), and a function can be declared with an operator_symbol as its name (see 6.1).
- 10 The syntax rules use the terms defining_identifier, defining_character_literal, and defining_operator_symbol for the defining occurrence of a name; these are collectively called *defining names*. The terms direct_name and selector_name are used for usage occurrences of identifiers, character_literals, and operator_symbols. These are collectively called *usage names*.

Dynamic Semantics

- 11 The process by which a construct achieves its run-time effect is called *execution*. This process is also called *elaboration* for declarations and *evaluation* for expressions. One of the terms execution, elaboration, or evaluation is defined by this International Standard for each construct that has a run-time effect.

NOTES

- 12 1 At compile time, the declaration of an entity *declares* the entity. At run time, the elaboration of the declaration *creates* the entity.

3.2 Types and Subtypes

Static Semantics

- 1 A *type* is characterized by a set of values, and a set of *primitive operations* which implement the fundamental aspects of its semantics. An *object* of a given type is a run-time entity that contains (has) a value of the type.
- 2/2 Types are grouped into *categories* of types. There exist several *language-defined categories* of types (see NOTES below), reflecting the similarity of their values and primitive operations. Most categories of types form *classes* of types. *Elementary* types are those whose values are logically indivisible; *composite* types are those whose values are composed of *component* values.
- 3 The elementary types are the *scalar* types (*discrete* and *real*) and the *access* types (whose values provide access to objects or subprograms). Discrete types are either *integer* types or are defined by enumeration of their values (*enumeration* types). Real types are either *floating point* types or *fixed point* types.
- 4/2 The composite types are the *record* types, *record extensions*, *array* types, *interface* types, *task* types, and *protected* types.
- 4.1/2 There can be multiple views of a type with varying sets of operations. An *incomplete* type represents an incomplete view (see 3.10.1) of a type with a very restricted usage, providing support for recursive data structures. A *private* type or *private extension* represents a partial view (see 7.3) of a type, providing support for data abstraction. The full view (see 3.2.1) of a type represents its complete definition. An incomplete or partial view is considered a composite type, even if the full view is not.
- 5/2 Certain composite types (and views thereof) have special components called *discriminants* whose values affect the presence, constraints, or initialization of other components. Discriminants can be thought of as parameters of the type.
- 6/2 The term *subcomponent* is used in this International Standard in place of the term component to indicate either a component, or a component of another subcomponent. Where other subcomponents are excluded, the term component is used instead. Similarly, a *part* of an object or value is used to mean the whole object or value, or any set of its subcomponents. The terms component, subcomponent, and part are also applied to a type meaning the component, subcomponent, or part of objects and values of the type.

The set of possible values for an object of a given type can be subjected to a condition that is called a *constraint* (the case of a *null constraint* that specifies no restriction is also included); the rules for which values satisfy a given kind of constraint are given in 3.5 for `range_constraints`, 3.6.1 for `index_constraints`, and 3.7.1 for `discriminant_constraints`. The set of possible values for an object of an access type can also be subjected to a condition that excludes the null value (see 3.10).

A *subtype* of a given type is a combination of the type, a constraint on values of the type, and certain attributes specific to the subtype. The given type is called the *type of the subtype*. Similarly, the associated constraint is called the *constraint of the subtype*. The set of values of a subtype consists of the values of its type that satisfy its constraint and any exclusion of the null value. Such values *belong* to the subtype.

A subtype is called an *unconstrained subtype* if its type has unknown discriminants, or if its type allows range, index, or discriminant constraints, but the subtype does not impose such a constraint; otherwise, the subtype is called a *constrained subtype* (since it has no unconstrained characteristics).

NOTES

2 Any set of types can be called a “category” of types, and any set of types that is closed under derivation (see 3.4) can be called a “class” of types. However, only certain categories and classes are used in the description of the rules of the language — generally those that have their own particular set of primitive operations (see 3.2.3), or that correspond to a set of types that are matched by a given kind of generic formal type (see 12.5). The following are examples of “interesting” *language-defined classes*: elementary, scalar, discrete, enumeration, character, boolean, integer, signed integer, modular, real, floating point, fixed point, ordinary fixed point, decimal fixed point, numeric, access, access-to-object, access-to-subprogram, composite, array, string, (untagged) record, tagged, task, protected, nonlimited. Special syntax is provided to define types in each of these classes. In addition to these classes, the following are examples of “interesting” *language-defined categories*: abstract, incomplete, interface, limited, private, record.

These language-defined categories are organized like this:

```

all types
    elementary
        scalar
            discrete
                enumeration
                    character
                    boolean
                    other enumeration
                integer
                    signed integer
                    modular integer
            real
                floating point
                fixed point
                    ordinary fixed point
                    decimal fixed point
        access
            access-to-object
            access-to-subprogram
    composite
        untagged
            array
                string
                other array
            record
            task
            protected
        tagged (including interfaces)
            nonlimited tagged record
            limited tagged
                limited tagged record

```

synchronized tagged
tagged task
tagged protected

- 13/2 There are other categories, such as “numeric” and “discriminated”, which represent other categorization dimensions, but do not fit into the above strictly hierarchical picture.

3.2.1 Type Declarations

- 1 A type_declaration declares a type and its first subtype.

Syntax

```

2   type_declaration ::= full_type_declaration
      | incomplete_type_declaration
      | private_type_declaration
      | private_extension_declaration

3   full_type_declaration ::= type defining_identifier [known_discriminant_part] is type_definition;
      | task_type_declaration
      | protected_type_declaration

4/2  type_definition ::= enumeration_type_definition | integer_type_definition
      | real_type_definition | array_type_definition
      | record_type_definition | access_type_definition
      | derived_type_definition | interface_type_definition

```

Legality Rules

- 5 A given type shall not have a subcomponent whose type is the given type itself.

Static Semantics

- 6 The defining_identifier of a type_declaration denotes the *first subtype* of the type. The known_discriminant_part, if any, defines the discriminants of the type (see 3.7, “Discriminants”). The remainder of the type_declaration defines the remaining characteristics of (the view of) the type.
- 7/2 A type defined by a type_declaration is a *named* type; such a type has one or more nameable subtypes. Certain other forms of declaration also include type definitions as part of the declaration for an object. The type defined by such a declaration is *anonymous* — it has no nameable subtypes. For explanatory purposes, this International Standard sometimes refers to an anonymous type by a pseudo-name, written in italics, and uses such pseudo-names at places where the syntax normally requires an identifier. For a named type whose first subtype is T, this International Standard sometimes refers to the type of T as simply “the type T”.
- 8/2 A named type that is declared by a full_type_declaration, or an anonymous type that is defined by an access_definition or as part of declaring an object of the type, is called a *full type*. The declaration of a full type also declares the *full view* of the type. The type_definition, task_definition, protected_definition, or access_definition that defines a full type is called a *full type definition*. Types declared by other forms of type_declaration are not separate types; they are partial or incomplete views of some full type.
- 9 The definition of a type implicitly declares certain *predefined operators* that operate on the type, according to what classes the type belongs, as specified in 4.5, “Operators and Expression Evaluation”.
- 10 The *predefined types* (for example the types Boolean, Wide_Character, Integer, root_integer, and universal_integer) are the types that are defined in a predefined library package called Standard; this

package also includes the (implicit) declarations of their predefined operators. The package Standard is described in A.1.

Dynamic Semantics

The elaboration of a **full_type_declaration** consists of the elaboration of the full type definition. Each elaboration of a full type definition creates a distinct type and its first subtype. 11

Examples

Examples of type definitions:

```
(White, Red, Yellow, Green, Blue, Brown, Black)
range 1 .. 72
array(1 .. 10) of Integer
```

Examples of type declarations:

```
type Color is (White, Red, Yellow, Green, Blue, Brown, Black);
type Column is range 1 .. 72;
type Table is array(1 .. 10) of Integer;
```

NOTES

3 Each of the above examples declares a named type. The identifier given denotes the first subtype of the type. Other named subtypes of the type can be declared with **subtype_declarations** (see 3.2.2). Although names do not directly denote types, a phrase like “the type Column” is sometimes used in this International Standard to refer to the type of Column, where Column denotes the first subtype of the type. For an example of the definition of an anonymous type, see the declaration of the array Color_Table in 3.3.1; its type is anonymous — it has no nameable subtypes. 16

3.2.2 Subtype Declarations

A **subtype_declaration** declares a subtype of some previously declared type, as defined by a **subtype_indication**. 1

Syntax

```
subtype_declaration ::= 2
  subtype_defining_identifier is subtype_indication;
subtype_indication ::= [null_exclusion] subtype_mark [constraint] 3/2
subtype_mark ::= subtype_name 4
constraint ::= scalar_constraint | composite_constraint 5
scalar_constraint ::= 6
  range_constraint | digits_constraint | delta_constraint
composite_constraint ::= 7
  index_constraint | discriminant_constraint
```

Name Resolution Rules

A **subtype_mark** shall resolve to denote a subtype. The type *determined by* a **subtype_mark** is the type of the subtype denoted by the **subtype_mark**. 8

Dynamic Semantics

The elaboration of a **subtype_declaration** consists of the elaboration of the **subtype_indication**. The elaboration of a **subtype_indication** creates a new subtype. If the **subtype_indication** does not include a **constraint**, the new subtype has the same (possibly null) constraint as that denoted by the **subtype_mark**. The elaboration of a **subtype_indication** that includes a **constraint** proceeds as follows: 9

- 10 • The constraint is first elaborated.
- 11 • A check is then made that the constraint is *compatible* with the subtype denoted by the subtype_mark.
- 12 The condition imposed by a constraint is the condition obtained after elaboration of the constraint. The rules defining compatibility are given for each form of constraint in the appropriate subclause. These rules are such that if a constraint is *compatible* with a subtype, then the condition imposed by the constraint cannot contradict any condition already imposed by the subtype on its values. The exception Constraint_Error is raised if any check of compatibility fails.

NOTES

13 4 A scalar_constraint may be applied to a subtype of an appropriate scalar type (see 3.5, 3.5.9, and J.3), even if the subtype is already constrained. On the other hand, a composite_constraint may be applied to a composite subtype (or an access-to-composite subtype) only if the composite subtype is unconstrained (see 3.6.1 and 3.7.1).

Examples

- 14 *Examples of subtype declarations:*

<pre>subtype Rainbow is Color range Red .. Blue; subtype Red_Blue is Rainbow; subtype Int is Integer; subtype Small_Int is Integer range -10 .. 10; subtype Up_To_K is Column range 1 .. K; subtype Square is Matrix(1 .. 10, 1 .. 10); subtype Male is Person(Sex => M); subtype Binop_Ref is not null Binop_Ptr;</pre>	<pre>-- see 3.2.1 -- see 3.2.1 -- see 3.6 -- see 3.10.1 -- see 3.10</pre>
--	---

3.2.3 Classification of Operations

Static Semantics

- 1/2 An operation *operates on a type T* if it yields a value of type T, if it has an operand whose expected type (see 8.6) is T, or if it has an access parameter or access result type (see 6.1) designating T. A predefined operator, or other language-defined operation such as assignment or a membership test, that operates on a type, is called a *predefined operation* of the type. The *primitive operations* of a type are the predefined operations of the type, plus any user-defined primitive subprograms.
- 2 The *primitive subprograms* of a specific type are defined as follows:
- 3 • The predefined operators of the type (see 4.5);
 - 4 • For a derived type, the inherited (see 3.4) user-defined subprograms;
 - 5 • For an enumeration type, the enumeration literals (which are considered parameterless functions — see 3.5.1);
 - 6 • For a specific type declared immediately within a package_specification, any subprograms (in addition to the enumeration literals) that are explicitly declared immediately within the same package_specification and that operate on the type;
- 7/2 • For a nonformal type, any subprograms not covered above that are explicitly declared immediately within the same declarative region as the type and that override (see 8.3) other implicitly declared primitive subprograms of the type.
- 8 A primitive subprogram whose designator is an operator_symbol is called a *primitive operator*.

3.3 Objects and Named Numbers

Objects are created at run time and contain a value of a given type. An object can be created and initialized as part of elaborating a declaration, evaluating an `allocator`, `aggregate`, or `function_call`, or passing a parameter by copy. Prior to reclaiming the storage for an object, it is finalized if necessary (see 7.6.1).

Static Semantics

All of the following are objects:

- the entity declared by an `object_declaration`;
- a formal parameter of a subprogram, entry, or generic subprogram;
- a generic formal object;
- a loop parameter;
- a choice parameter of an `exception_handler`;
- an entry index of an `entry_body`;
- the result of dereferencing an access-to-object value (see 4.1);
- the return object created as the result of evaluating a `function_call` (or the equivalent operator invocation — see 6.6);
- the result of evaluating an `aggregate`;
- a component, slice, or view conversion of another object.

An object is either a *constant* object or a *variable* object. The value of a constant object cannot be changed between its initialization and its finalization, whereas the value of a variable object can be changed. Similarly, a view of an object is either a *constant* or a *variable*. All views of a constant object are constant. A constant view of a variable object cannot be used to modify the value of the variable. The terms *constant* and *variable* by themselves refer to constant and variable views of objects.

The value of an object is *read* when the value of any part of the object is evaluated, or when the value of an enclosing object is evaluated. The value of a variable is *updated* when an assignment is performed to any part of the variable, or when an assignment is performed to an enclosing object.

Whether a view of an object is constant or variable is determined by the definition of the view. The following (and no others) represent constants:

- an object declared by an `object_declaration` with the reserved word **constant**;
- a formal parameter or generic formal object of mode **in**;
- a discriminant;
- a loop parameter, choice parameter, or entry index;
- the dereference of an access-to-constant value;
- the result of evaluating a `function_call` or an `aggregate`;
- a `selected_component`, `indexed_component`, slice, or view conversion of a constant.

At the place where a view of an object is defined, a *nominal subtype* is associated with the view. The object's *actual subtype* (that is, its subtype) can be more restrictive than the nominal subtype of the view; it always is if the nominal subtype is an *indefinite subtype*. A subtype is an indefinite subtype if it is an unconstrained array subtype, or if it has unknown discriminants or unconstrained discriminants without

defaults (see 3.7); otherwise the subtype is a *definite* subtype (all elementary subtypes are definite subtypes). A class-wide subtype is defined to have unknown discriminants, and is therefore an indefinite subtype. An indefinite subtype does not by itself provide enough information to create an object; an additional constraint or explicit initialization expression is necessary (see 3.3.1). A component cannot have an indefinite nominal subtype.

- 24 A *named number* provides a name for a numeric value known at compile time. It is declared by a `number_declaration`.

NOTES

- 25 5 A constant cannot be the target of an assignment operation, nor be passed as an `in out` or `out` parameter, between its initialization and finalization, if any.
- 26 6 The nominal and actual subtypes of an elementary object are always the same. For a discriminated or array object, if the nominal subtype is constrained then so is the actual subtype.

3.3.1 Object Declarations

- 1 An `object_declaration` declares a *stand-alone* object with a given nominal subtype and, optionally, an explicit initial value given by an initialization expression. For an array, task, or protected object, the `object_declaration` may include the definition of the (anonymous) type of the object.

Syntax

- 2/2 `object_declaration ::=`
 `defining_identifier_list : [aliased] [constant] subtype_indication [:= expression];`
 `| defining_identifier_list : [aliased] [constant] access_definition [:= expression];`
 `| defining_identifier_list : [aliased] [constant] array_type_definition [:= expression];`
 `| single_task_declaration`
 `| single_protected_declaration`
- 3 `defining_identifier_list ::=`
 `defining_identifier {, defining_identifier}`

Name Resolution Rules

- 4 For an `object_declaration` with an `expression` following the compound delimiter `:=`, the type expected for the `expression` is that of the object. This `expression` is called the *initialization expression*.

Legality Rules

- 5/2 An `object_declaration` without the reserved word `constant` declares a variable object. If it has a `subtype_indication` or an `array_type_definition` that defines an indefinite subtype, then there shall be an initialization expression.

Static Semantics

- 6 An `object_declaration` with the reserved word `constant` declares a constant object. If it has an initialization expression, then it is called a *full constant declaration*. Otherwise it is called a *deferred constant declaration*. The rules for deferred constant declarations are given in clause 7.4. The rules for full constant declarations are given in this subclause.
- 7 Any declaration that includes a `defining_identifier_list` with more than one `defining_identifier` is equivalent to a series of declarations each containing one `defining_identifier` from the list, with the rest of the text of the declaration copied for each declaration in the series, in the same order as the list. The remainder of this International Standard relies on this equivalence; explanations are given for declarations with a single `defining_identifier`.

The `subtype_indication`, `access_definition`, or full type definition of an `object_declaration` defines the nominal subtype of the object. The `object_declaration` declares an object of the type of the nominal subtype.

A component of an object is said to *require late initialization* if it has an access discriminant value constrained by a per-object expression, or if it has an initialization expression that includes a name denoting the current instance of the type or denoting an access discriminant.

Dynamic Semantics

If a composite object declared by an `object_declaration` has an unconstrained nominal subtype, then if this subtype is indefinite or the object is constant the actual subtype of this object is constrained. The constraint is determined by the bounds or discriminants (if any) of its initial value; the object is said to be *constrained by its initial value*. When not constrained by its initial value, the actual and nominal subtypes of the object are the same. If its actual subtype is constrained, the object is called a *constrained object*.

For an `object_declaration` without an initialization expression, any initial values for the object or its subcomponents are determined by the *implicit initial values* defined for its nominal subtype, as follows:

- The implicit initial value for an access subtype is the null value of the access type.
- The implicit initial (and only) value for each discriminant of a constrained discriminated subtype is defined by the subtype.
- For a (definite) composite subtype, the implicit initial value of each component with a `default_expression` is obtained by evaluation of this expression and conversion to the component's nominal subtype (which might raise `Constraint_Error` — see 4.6, “Type Conversions”), unless the component is a discriminant of a constrained subtype (the previous case), or is in an excluded variant (see 3.8.1). For each component that does not have a `default_expression`, any implicit initial values are those determined by the component's nominal subtype.
- For a protected or task subtype, there is an implicit component (an entry queue) corresponding to each entry, with its implicit initial value being an empty queue.

The elaboration of an `object_declaration` proceeds in the following sequence of steps:

1. The `subtype_indication`, `access_definition`, `array_type_definition`, `single_task_declaration`, or `single_protected_declaration` is first elaborated. This creates the nominal subtype (and the anonymous type in the last four cases).
2. If the `object_declaration` includes an initialization expression, the (explicit) initial value is obtained by evaluating the expression and converting it to the nominal subtype (which might raise `Constraint_Error` — see 4.6).
3. The object is created, and, if there is not an initialization expression, the object is *initialized by default*. When an object is initialized by default, any per-object constraints (see 3.8) are elaborated and any implicit initial values for the object or for its subcomponents are obtained as determined by the nominal subtype. Any initial values (whether explicit or implicit) are assigned to the object or to the corresponding subcomponents. As described in 5.2 and 7.6, Initialize and Adjust procedures can be called.

This paragraph was deleted.

For the third step above, evaluations and assignments are performed in an arbitrary order subject to the following restrictions:

- Assignment to any part of the object is preceded by the evaluation of the value that is to be assigned.

- 20.2/2 • The evaluation of a `default_expression` that includes the name of a discriminant is preceded by the assignment to that discriminant.
- 20.3/2 • The evaluation of the `default_expression` for any component that depends on a discriminant is preceded by the assignment to that discriminant.
- 20.4/2 • The assignments to any components, including implicit components, not requiring late initialization must precede the initial value evaluations for any components requiring late initialization; if two components both require late initialization, then assignments to parts of the component occurring earlier in the order of the component declarations must precede the initial value evaluations of the component occurring later.
- 21 There is no implicit initial value defined for a scalar subtype. In the absence of an explicit initialization, a newly created scalar object might have a value that does not belong to its subtype (see 13.9.1 and H.1).
- 22 NOTES
- 22 7 Implicit initial values are not defined for an indefinite subtype, because if an object's nominal subtype is indefinite, an explicit initial value is required.
- 23 8 As indicated above, a stand-alone object is an object declared by an `object_declaration`. Similar definitions apply to “stand-alone constant” and “stand-alone variable.” A subcomponent of an object is not a stand-alone object, nor is an object that is created by an `allocator`. An object declared by a `loop_parameter_specification`, `parameter_specification`, `entry_index_specification`, `choice_parameter_specification`, or a `formal_object_declaration` is not called a stand-alone object.
- 24 9 The type of a stand-alone object cannot be abstract (see 3.9.3).

Examples

25 *Example of a multiple object declaration:*

26 *-- the multiple object declaration*

27/2 `John, Paul : not null Person_Name := new Person(Sex => M); -- see 3.10.1`

28 *-- is equivalent to the two single object declarations in the order given*

29/2 `John : not null Person_Name := new Person(Sex => M);
Paul : not null Person_Name := new Person(Sex => M);`

30 *Examples of variable declarations:*

31/2 `Count, Sum : Integer;
Size : Integer range 0 .. 10_000 := 0;
Sorted : Boolean := False;
Color_Table : array(1 .. Max) of Color;
Option : Bit_Vector(1 .. 10) := (others => True);
Hello : aliased String := "Hi, world.";
θ, φ : Float range -π .. +π;`

32 *Examples of constant declarations:*

33/2 `Limit : constant Integer := 10_000;
Low_Limit : constant Integer := Limit/10;
Tolerance : constant Real := Dispersion(1.15);
Hello_Msg : constant access String := Hello'Access; -- see 3.10.2`

3.3.2 Number Declarations

1 A `number_declaration` declares a named number.

Syntax

2 `number_declaration ::=
defining_identifier_list : constant := static_expression;`

Name Resolution Rules

The *static_expression* given for a *number_declaration* is expected to be of any numeric type.

3

Legality Rules

The *static_expression* given for a number declaration shall be a static expression, as defined by clause 4.9.

4

Static Semantics

The named number denotes a value of type *universal_integer* if the type of the *static_expression* is an integer type. The named number denotes a value of type *universal_real* if the type of the *static_expression* is a real type.

5

The value denoted by the named number is the value of the *static_expression*, converted to the corresponding universal type.

6

Dynamic Semantics

The elaboration of a *number_declaration* has no effect.

7

*Examples**Examples of number declarations:*

Two_Pi	<code>: constant := 2.0*Ada.Numerics.Pi;</code>	-- a real number (see A.5)
Max	<code>: constant := 500;</code>	-- an integer number
Max_Line_Size	<code>: constant := Max/6;</code>	-- the integer 83
Power_16	<code>: constant := 2**16;</code>	-- the integer 65_536
One, Ün, Eins	<code>: constant := 1;</code>	-- three different names for 1

8

9

10/2

3.4 Derived Types and Classes

A *derived_type_definition* defines a *derived type* (and its first subtype) whose characteristics are *derived* from those of a parent type, and possibly from progenitor types.

1/2

A *class of types* is a set of types that is closed under derivation; that is, if the parent or a progenitor type of a derived type belongs to a class, then so does the derived type. By saying that a particular group of types forms a class, we are saying that all derivatives of a type in the set inherit the characteristics that define that set. The more general term *category of types* is used for a set of types whose defining characteristics are not necessarily inherited by derivatives; for example, limited, abstract, and interface are all categories of types, but not classes of types.

1.1/2

Syntax

```
derived_type_definition ::= [abstract] [limited] new parent_subtype_indication [[and interface_list] record_extension_part]
```

2/2

Legality Rules

The *parent_subtype_indication* defines the *parent subtype*; its type is the *parent type*. The *interface_list* defines the progenitor types (see 3.9.4). A derived type has one parent type and zero or more progenitor types.

3/2

A type shall be completely defined (see 3.11.1) prior to being specified as the parent type in a *derived_type_definition* — the *full_type_declarations* for the parent type and any of its subcomponents have to precede the *derived_type_definition*.

4

- 5/2 If there is a `record_extension_part`, the derived type is called a *record extension* of the parent type. A `record_extension_part` shall be provided if and only if the parent type is a tagged type. An `interface_list` shall be provided only if the parent type is a tagged type.
- 5.1/2 If the reserved word **limited** appears in a `derived_type_definition`, the parent type shall be a limited type.

Static Semantics

- 6 The first subtype of the derived type is unconstrained if a `known_discriminant_part` is provided in the declaration of the derived type, or if the parent subtype is unconstrained. Otherwise, the constraint of the first subtype *corresponds* to that of the parent subtype in the following sense: it is the same as that of the parent subtype except that for a range constraint (implicit or explicit), the value of each bound of its range is replaced by the corresponding value of the derived type.
- 6.1/2 The first subtype of the derived type excludes null (see 3.10) if and only if the parent subtype excludes null.
- 7 The characteristics of the derived type are defined as follows:
- 8/2
- If the parent type or a progenitor type belongs to a class of types, then the derived type also belongs to that class. The following sets of types, as well as any higher-level sets composed from them, are classes in this sense, and hence the characteristics defining these classes are inherited by derived types from their parent or progenitor types: signed integer, modular integer, ordinary fixed, decimal fixed, floating point, enumeration, boolean, character, access-to-constant, general access-to-variable, pool-specific access-to-variable, access-to-subprogram, array, string, non-array composite, nonlimited, untagged record, tagged, task, protected, and synchronized tagged.
- 9
- If the parent type is an elementary type or an array type, then the set of possible values of the derived type is a copy of the set of possible values of the parent type. For a scalar type, the base range of the derived type is the same as that of the parent type.
- 10
- If the parent type is a composite type other than an array type, then the components, protected subprograms, and entries that are declared for the derived type are as follows:
 - 11 • The discriminants specified by a new `known_discriminant_part`, if there is one; otherwise, each discriminant of the parent type (implicitly declared in the same order with the same specifications) — in the latter case, the discriminants are said to be *inherited*, or if unknown in the parent, are also unknown in the derived type;
 - 12 • Each nondiscriminant component, entry, and protected subprogram of the parent type, implicitly declared in the same order with the same declarations; these components, entries, and protected subprograms are said to be *inherited*;
 - 13 • Each component declared in a `record_extension_part`, if any.
- 14 Declarations of components, protected subprograms, and entries, whether implicit or explicit, occur immediately within the declarative region of the type, in the order indicated above, following the parent `subtype_indication`.
- 15/2
- *This paragraph was deleted.*
- 16
- For each predefined operator of the parent type, there is a corresponding predefined operator of the derived type.
- 17/2
- For each user-defined primitive subprogram (other than a user-defined equality operator — see below) of the parent type or of a progenitor type that already exists at the place of the `derived_type_definition`, there exists a corresponding *inherited* primitive subprogram of the derived type with the same defining name. Primitive user-defined equality operators of the

parent type and any progenitor types are also inherited by the derived type, except when the derived type is a nonlimited record extension, and the inherited operator would have a profile that is type conformant with the profile of the corresponding predefined equality operator; in this case, the user-defined equality operator is not inherited, but is rather incorporated into the implementation of the predefined equality operator of the record extension (see 4.5.2).

The profile of an inherited subprogram (including an inherited enumeration literal) is obtained from the profile of the corresponding (user-defined) primitive subprogram of the parent or progenitor type, after systematic replacement of each subtype of its profile (see 6.1) that is of the parent or progenitor type with a *corresponding subtype* of the derived type. For a given subtype of the parent or progenitor type, the corresponding subtype of the derived type is defined as follows:

- If the declaration of the derived type has neither a `known_discriminant_part` nor a `record_extension_part`, then the corresponding subtype has a constraint that corresponds (as defined above for the first subtype of the derived type) to that of the given subtype.
- If the derived type is a record extension, then the corresponding subtype is the first subtype of the derived type.
- If the derived type has a new `known_discriminant_part` but is not a record extension, then the corresponding subtype is constrained to those values that when converted to the parent type belong to the given subtype (see 4.6).

The same formal parameters have `default_expressions` in the profile of the inherited subprogram. Any type mismatch due to the systematic replacement of the parent or progenitor type by the derived type is handled as part of the normal type conversion associated with parameter passing — see 6.4.1.

If a primitive subprogram of the parent or progenitor type is visible at the place of the `derived_type_definition`, then the corresponding inherited subprogram is implicitly declared immediately after the `derived_type_definition`. Otherwise, the inherited subprogram is implicitly declared later or not at all, as explained in 7.3.1.

A derived type can also be defined by a `private_extension_declaration` (see 7.3) or a `formal_derived_type_definition` (see 12.5.1). Such a derived type is a partial view of the corresponding full or actual type.

All numeric types are derived types, in that they are implicitly derived from a corresponding root numeric type (see 3.5.4 and 3.5.6).

Dynamic Semantics

The elaboration of a `derived_type_definition` creates the derived type and its first subtype, and consists of the elaboration of the `subtype_indication` and the `record_extension_part`, if any. If the `subtype_indication` depends on a discriminant, then only those expressions that do not depend on a discriminant are evaluated.

For the execution of a call on an inherited subprogram, a call on the corresponding primitive subprogram of the parent or progenitor type is performed; the normal conversion of each actual parameter to the subtype of the corresponding formal parameter (see 6.4.1) performs any necessary type conversion as well. If the result type of the inherited subprogram is the derived type, the result of calling the subprogram of the parent or progenitor is converted to the derived type, or in the case of a null extension, extended to the derived type using the equivalent of an `extension_aggregate` with the original result as the `ancestor_part` and `null_record` as the `record_component_association_list`.

NOTES

- 28 10 Classes are closed under derivation — any class that contains a type also contains its derivatives. Operations available for a given class of types are available for the derived types in that class.
- 29 11 Evaluating an inherited enumeration literal is equivalent to evaluating the corresponding enumeration literal of the parent type, and then converting the result to the derived type. This follows from their equivalence to parameterless functions.
- 30 12 A generic subprogram is not a subprogram, and hence cannot be a primitive subprogram and cannot be inherited by a derived type. On the other hand, an instance of a generic subprogram can be a primitive subprogram, and hence can be inherited.
- 31 13 If the parent type is an access type, then the parent and the derived type share the same storage pool; there is a **null** access value for the derived type and it is the implicit initial value for the type. See 3.10.
- 32 14 If the parent type is a boolean type, the predefined relational operators of the derived type deliver a result of the predefined type Boolean (see 4.5.2). If the parent type is an integer type, the right operand of the predefined exponentiation operator is of the predefined type Integer (see 4.5.6).
- 33 15 Any discriminants of the parent type are either all inherited, or completely replaced with a new set of discriminants.
- 34 16 For an inherited subprogram, the subtype of a formal parameter of the derived type need not have any value in common with the first subtype of the derived type.
- 35 17 If the reserved word **abstract** is given in the declaration of a type, the type is abstract (see 3.9.3).
- 35.1/2 18 An interface type that has a progenitor type “is derived from” that type. A **derived_type_definition**, however, never defines an interface type.
- 35.2/2 19 It is illegal for the parent type of a **derived_type_definition** to be a synchronized tagged type.

Examples

36 Examples of derived type declarations:

```
37    type Local_Coordinate is new Coordinate;      -- two different types
      type Midweek is new Day range Tue .. Thu;    -- see 3.5.1
      type Counter is new Positive;                 -- same range as Positive
38    type Special_Key is new Key_Manager.Key;     -- see 7.3.1
      -- the inherited subprograms have the following specifications:
      -- procedure Get_Key(K : out Special_Key);
      -- function "<"(X,Y : Special_Key) return Boolean;
```

3.4.1 Derivation Classes

- 1 In addition to the various language-defined classes of types, types can be grouped into *derivation classes*.

Static Semantics

- 2/2 A derived type is *derived from* its parent type *directly*; it is derived *indirectly* from any type from which its parent type is derived. A derived type, interface type, type extension, task type, protected type, or formal derived type is also derived from every ancestor of each of its progenitor types, if any. The derivation class of types for a type *T* (also called the class *rooted at T*) is the set consisting of *T* (the *root type* of the class) and all types derived from *T* (directly or indirectly) plus any associated universal or class-wide types (defined below).
- 3/2 Every type is either a *specific type*, a *class-wide type*, or a *universal type*. A specific type is one defined by a **type_declaration**, a **formal_type_declaration**, or a full type definition embedded in another construct. Class-wide and universal types are implicitly defined, to act as representatives for an entire class of types, as follows:

Class-wide types

4
Class-wide types are defined for (and belong to) each derivation class rooted at a tagged type (see 3.9). Given a subtype S of a tagged type T, S'Class is the *subtype_mark* for a corresponding subtype of the tagged class-wide type T'Class. Such types are called “class-wide” because when a formal parameter is defined to be of a class-wide type T'Class, an actual parameter of any type in the derivation class rooted at T is acceptable (see 8.6).

5
The set of values for a class-wide type T'Class is the discriminated union of the set of values of each specific type in the derivation class rooted at T (the tag acts as the implicit discriminant — see 3.9). Class-wide types have no primitive subprograms of their own. However, as explained in 3.9.2, operands of a class-wide type T'Class can be used as part of a dispatching call on a primitive subprogram of the type T. The only components (including discriminants) of T'Class that are visible are those of T. If S is a first subtype, then S'Class is a first subtype.

Universal types

6/2
Universal types are defined for (and belong to) the integer, real, fixed point, and access classes, and are referred to in this standard as respectively, *universal_integer*, *universal_real*, *universal_fixed*, and *universal_access*. These are analogous to class-wide types for these language-defined elementary classes. As with class-wide types, if a formal parameter is of a universal type, then an actual parameter of any type in the corresponding class is acceptable. In addition, a value of a universal type (including an integer or real *numeric_literal*, or the literal **null**) is “universal” in that it is acceptable where some particular type in the class is expected (see 8.6).

7
The set of values of a universal type is the undiscriminated union of the set of values possible for any definable type in the associated class. Like class-wide types, universal types have no primitive subprograms of their own. However, their “universality” allows them to be used as operands with the primitive subprograms of any type in the corresponding class.

8
The integer and real numeric classes each have a specific root type in addition to their universal type, named respectively *root_integer* and *root_real*.

9
A class-wide or universal type is said to *cover* all of the types in its class. A specific type covers only itself.

10/2
A specific type T2 is defined to be a *descendant* of a type T1 if T2 is the same as T1, or if T2 is derived (directly or indirectly) from T1. A class-wide type T2'Class is defined to be a descendant of type T1 if T2 is a descendant of T1. Similarly, the numeric universal types are defined to be descendants of the root types of their classes. If a type T2 is a descendant of a type T1, then T1 is called an *ancestor* of T2. An *ultimate ancestor* of a type is an ancestor of that type that is not itself a descendant of any other type. Every untagged type has a unique ultimate ancestor.

11
An inherited component (including an inherited discriminant) of a derived type is inherited *from* a given ancestor of the type if the corresponding component was inherited by each derived type in the chain of derivations going back to the given ancestor.

NOTES

12
20 Because operands of a universal type are acceptable to the predefined operators of any type in their class, ambiguity can result. For *universal_integer* and *universal_real*, this potential ambiguity is resolved by giving a preference (see 8.6) to the predefined operators of the corresponding root types (*root_integer* and *root_real*, respectively). Hence, in an apparently ambiguous expression like

13
 $1 + 4 < 7$

14
where each of the literals is of type *universal_integer*, the predefined operators of *root_integer* will be preferred over those of other specific integer types, thereby resolving the ambiguity.

3.5 Scalar Types

- 1 *Scalar* types comprise enumeration types, integer types, and real types. Enumeration types and integer types are called *discrete* types; each value of a discrete type has a *position number* which is an integer value. Integer types and real types are called *numeric* types. All scalar types are ordered, that is, all relational operators are predefined for their values.

Syntax

```
2     range_constraint ::= range range
3     range ::= range_attribute_reference
| simple_expression .. simple_expression
```

- 4 A *range* has a *lower bound* and an *upper bound* and specifies a subset of the values of some scalar type (the *type of the range*). A range with lower bound L and upper bound R is described by “L .. R”. If R is less than L, then the range is a *null range*, and specifies an empty set of values. Otherwise, the range specifies the values of the type from the lower bound to the upper bound, inclusive. A value *belongs* to a range if it is of the type of the range, and is in the subset of values specified by the range. A value *satisfies* a range constraint if it belongs to the associated range. One range is *included* in another if all values that belong to the first range also belong to the second.

Name Resolution Rules

- 5 For a *subtype_indication* containing a *range_constraint*, either directly or as part of some other *scalar_constraint*, the type of the *range* shall resolve to that of the type determined by the *subtype_mark* of the *subtype_indication*. For a *range* of a given type, the *simple_expressions* of the *range* (likewise, the *simple_expressions* of the equivalent *range* for a *range_attribute_reference*) are expected to be of the type of the *range*.

Static Semantics

- 6 The *base range* of a scalar type is the range of finite values of the type that can be represented in every unconstrained object of the type; it is also the range supported at a minimum for intermediate values during the evaluation of expressions involving predefined operators of the type.
- 7 A constrained scalar subtype is one to which a range constraint applies. The *range* of a constrained scalar subtype is the range associated with the range constraint of the subtype. The *range* of an unconstrained scalar subtype is the base range of its type.

Dynamic Semantics

- 8 A *range* is *compatible* with a scalar subtype if and only if it is either a null range or each bound of the range belongs to the range of the subtype. A *range_constraint* is *compatible* with a scalar subtype if and only if its range is compatible with the subtype.
- 9 The elaboration of a *range_constraint* consists of the evaluation of the *range*. The evaluation of a *range* determines a lower bound and an upper bound. If *simple_expressions* are given to specify bounds, the evaluation of the *range* evaluates these *simple_expressions* in an arbitrary order, and converts them to the type of the *range*. If a *range_attribute_reference* is given, the evaluation of the *range* consists of the evaluation of the *range_attribute_reference*.

<i>Attributes</i>	10	
For every scalar subtype S, the following attributes are defined:	11	
S'First	S'First denotes the lower bound of the range of S. The value of this attribute is of the type of S.	12
S'Last	S'Last denotes the upper bound of the range of S. The value of this attribute is of the type of S.	13
S'Range	S'Range is equivalent to the range S'First .. S'Last.	14
S'Base	S'Base denotes an unconstrained subtype of the type of S. This unconstrained subtype is called the <i>base subtype</i> of the type.	15
S'Min	S'Min denotes a function with the following specification: <pre>function S'Min(Left, Right : S'Base) return S'Base</pre>	16
	The function returns the lesser of the values of the two parameters.	17
S'Max	S'Max denotes a function with the following specification: <pre>function S'Max(Left, Right : S'Base) return S'Base</pre>	18
	The function returns the greater of the values of the two parameters.	19
S'Succ	S'Succ denotes a function with the following specification: <pre>function S'Succ(Arg : S'Base) return S'Base</pre>	20
	For an enumeration type, the function returns the value whose position number is one more than that of the value of Arg; Constraint_Error is raised if there is no such value of the type. For an integer type, the function returns the result of adding one to the value of Arg. For a fixed point type, the function returns the result of adding small to the value of Arg. For a floating point type, the function returns the machine number (as defined in 3.5.7) immediately above the value of Arg; Constraint_Error is raised if there is no such machine number.	21
S'Pred	S'Pred denotes a function with the following specification: <pre>function S'Pred(Arg : S'Base) return S'Base</pre>	22
	For an enumeration type, the function returns the value whose position number is one less than that of the value of Arg; Constraint_Error is raised if there is no such value of the type. For an integer type, the function returns the result of subtracting one from the value of Arg. For a fixed point type, the function returns the result of subtracting small from the value of Arg. For a floating point type, the function returns the machine number (as defined in 3.5.7) immediately below the value of Arg; Constraint_Error is raised if there is no such machine number.	23
S'Wide_Wide_Image	S'Wide_Wide_Image denotes a function with the following specification: <pre>function S'Wide_Wide_Image(Arg : S'Base) return Wide_Wide_String</pre>	24
	The function returns an <i>image</i> of the value of Arg, that is, a sequence of characters representing the value in display form. The lower bound of the result is one.	25
	The image of an integer value is the corresponding decimal literal, without underlines, leading zeros, exponent, or trailing spaces, but with a single leading character that is either a minus sign or a space.	26
		27.1/2
		27.2/2
		27.3/2
		27.4/2

27.5/2

The image of an enumeration value is either the corresponding identifier in upper case or the corresponding character literal (including the two apostrophes); neither leading nor trailing spaces are included. For a *nongraphic character* (a value of a character type that has no enumeration literal associated with it), the result is a corresponding language-defined name in upper case (for example, the image of the nongraphic character identified as *nul* is “NUL” — the quotes are not part of the image).

27.6/2

The image of a floating point value is a decimal real literal best approximating the value (rounded away from zero if halfway between) with a single leading character that is either a minus sign or a space, a single digit (that is nonzero unless the value is zero), a decimal point, S'Digits-1 (see 3.5.8) digits after the decimal point (but one if S'Digits is one), an upper case E, the sign of the exponent (either + or -), and two or more digits (with leading zeros if necessary) representing the exponent. If S'Signed_Zeros is True, then the leading character is a minus sign for a negatively signed zero.

27.7/2

The image of a fixed point value is a decimal real literal best approximating the value (rounded away from zero if halfway between) with a single leading character that is either a minus sign or a space, one or more digits before the decimal point (with no redundant leading zeros), a decimal point, and S'Aft (see 3.5.10) digits after the decimal point.

28 S'Wide_Image S'Wide_Image denotes a function with the following specification:

```
29      function S'Wide_Image(Arg : S'Base)
           return Wide_String
```

30/2

The function returns an image of the value of Arg as a Wide_String. The lower bound of the result is one. The image has the same sequence of character as defined for S'Wide_Wide_Image if all the graphic characters are defined in Wide_Character; otherwise the sequence of characters is implementation defined (but no shorter than that of S'Wide_Wide_Image for the same value of Arg).

Paragraphs 31 through 34 were moved to Wide_Wide_Image.

35 S'Image S'Image denotes a function with the following specification:

```
36      function S'Image(Arg : S'Base)
           return String
```

37/2

The function returns an image of the value of Arg as a String. The lower bound of the result is one. The image has the same sequence of graphic characters as that defined for S'Wide_Wide_Image if all the graphic characters are defined in Character; otherwise the sequence of characters is implementation defined (but no shorter than that of S'Wide_Wide_Image for the same value of Arg).

37.1/2 S'Wide_Wide_Width

S'Wide_Wide_Width denotes the maximum length of a Wide_Wide_String returned by S'Wide_Wide_Image over all values of the subtype S. It denotes zero for a subtype that has a null range. Its type is *universal_integer*.

38 S'Wide_Width

S'Wide_Width denotes the maximum length of a Wide_String returned by S'Wide_Image over all values of the subtype S. It denotes zero for a subtype that has a null range. Its type is *universal_integer*.

39 S'Width

S'Width denotes the maximum length of a String returned by S'Image over all values of the subtype S. It denotes zero for a subtype that has a null range. Its type is *universal_integer*.

39.1/2 S'Wide_Wide_Value

S'Wide_Wide_Value denotes a function with the following specification:

```
39.2/2      function S'Wide_Wide_Value(Arg : Wide_Wide_String)
                  return S'Base
```

This function returns a value given an image of the value as a `Wide_Wide_String`, ignoring any leading or trailing spaces. 39.3/2

For the evaluation of a call on `S'Wide_Wide_Value` for an enumeration subtype `S`, if the sequence of characters of the parameter (ignoring leading and trailing spaces) has the syntax of an enumeration literal and if it corresponds to a literal of the type of `S` (or corresponds to the result of `S'Wide_Wide_Image` for a nongraphic character of the type), the result is the corresponding enumeration value; otherwise `Constraint_Error` is raised. 39.4/2

For the evaluation of a call on `S'Wide_Wide_Value` for an integer subtype `S`, if the sequence of characters of the parameter (ignoring leading and trailing spaces) has the syntax of an integer literal, with an optional leading sign character (plus or minus for a signed type; only plus for a modular type), and the corresponding numeric value belongs to the base range of the type of `S`, then that value is the result; otherwise `Constraint_Error` is raised. 39.5/2

For the evaluation of a call on `S'Wide_Wide_Value` for a real subtype `S`, if the sequence of characters of the parameter (ignoring leading and trailing spaces) has the syntax of one of the following: 39.6/2

- `numeric_literal` 39.7/2
- `numeral.[exponent]` 39.8/2
- `.numeral[exponent]` 39.9/2
- `base#based_numeral.#[exponent]` 39.10/2
- `base#.based_numeral#[exponent]` 39.11/2

with an optional leading sign character (plus or minus), and if the corresponding numeric value belongs to the base range of the type of `S`, then that value is the result; otherwise `Constraint_Error` is raised. The sign of a zero value is preserved (positive if none has been specified) if `S'Signed_Zeros` is `True`. 39.12/2

`S'Wide_Value` 40

`S'Wide_Value` denotes a function with the following specification:

```
function S'Wide_Value(Arg : Wide_String)
  return S'Base
```

This function returns a value given an image of the value as a `Wide_String`, ignoring any leading or trailing spaces. 42

For the evaluation of a call on `S'Wide_Value` for an enumeration subtype `S`, if the sequence of characters of the parameter (ignoring leading and trailing spaces) has the syntax of an enumeration literal and if it corresponds to a literal of the type of `S` (or corresponds to the result of `S'Wide_Image` for a value of the type), the result is the corresponding enumeration value; otherwise `Constraint_Error` is raised. For a numeric subtype `S`, the evaluation of a call on `S'Wide_Value` with `Arg` of type `Wide_String` is equivalent to a call on `S'Wide_Wide_Value` for a corresponding `Arg` of type `Wide_Wide_String`. 43/2

Paragraphs 44 through 51 were moved to `Wide_Wide_Value`.

`S'Value` 52

`S'Value` denotes a function with the following specification:

```
function S'Value(Arg : String)
  return S'Base
```

This function returns a value given an image of the value as a `String`, ignoring any leading or trailing spaces. 54

For the evaluation of a call on `S'Value` for an enumeration subtype `S`, if the sequence of characters of the parameter (ignoring leading and trailing spaces) has the syntax of an

enumeration literal and if it corresponds to a literal of the type of S (or corresponds to the result of S'Image for a value of the type), the result is the corresponding enumeration value; otherwise Constraint_Error is raised. For a numeric subtype S, the evaluation of a call on S'Value with Arg of type String is equivalent to a call on S'Wide_Wide_Value for a corresponding Arg of type Wide_Wide_String.

Implementation Permissions

- 56/2 An implementation may extend the Wide_Wide_Value, Wide_Value, Value, Wide_Wide_Image, Wide_Image, and Image attributes of a floating point type to support special values such as infinities and NaNs.

NOTES

- 57 21 The evaluation of S'First or S'Last never raises an exception. If a scalar subtype S has a nonnull range, S'First and S'Last belong to this range. These values can, for example, always be assigned to a variable of subtype S.
- 58 22 For a subtype of a scalar type, the result delivered by the attributes Succ, Pred, and Value might not belong to the subtype; similarly, the actual parameters of the attributes Succ, Pred, and Image need not belong to the subtype.
- 59 23 For any value V (including any nongraphic character) of an enumeration subtype S, S'Value(S'Image(V)) equals V, as do S'Wide_Value(S'Wide_Image(V)) and S'Wide_Wide_Value(S'Wide_Wide_Image(V)). None of these expressions ever raise Constraint_Error.

Examples

- 60 Examples of ranges:

```
-10 .. 10
X .. X + 1
0.0 .. 2.0*Pi
Red .. Green      -- see 3.5.1
1 .. 0           -- a null range
Table'Range       -- a range attribute reference (see 3.6)
```

- 62 Examples of range constraints:

```
range -999.0 .. +999.0
range S'First+1 .. S'Last-1
```

3.5.1 Enumeration Types

- 1 An enumeration_type_definition defines an enumeration type.

Syntax

- 2 enumeration_type_definition ::=
 (enumeration_literal_specification {, enumeration_literal_specification})
- 3 enumeration_literal_specification ::= defining_identifier | defining_character_literal
- 4 defining_character_literal ::= character_literal

Legality Rules

- 5 The defining_identifiers and defining_character_literals listed in an enumeration_type_definition shall be distinct.

Static Semantics

- 6 Each enumeration_literal_specification is the explicit declaration of the corresponding *enumeration literal*: it declares a parameterless function, whose defining name is the defining_identifier or defining_character_literal, and whose result type is the enumeration type.

Each enumeration literal corresponds to a distinct value of the enumeration type, and to a distinct position number. The position number of the value of the first listed enumeration literal is zero; the position number of the value of each subsequent enumeration literal is one more than that of its predecessor in the list.

The predefined order relations between values of the enumeration type follow the order of corresponding position numbers.

If the same `defining_identifier` or `defining_character_literal` is specified in more than one `enumeration_type_definition`, the corresponding enumeration literals are said to be *overloaded*. At any place where an overloaded enumeration literal occurs in the text of a program, the type of the enumeration literal has to be determinable from the context (see 8.6).

Dynamic Semantics

The elaboration of an `enumeration_type_definition` creates the enumeration type and its first subtype, which is constrained to the base range of the type.

When called, the parameterless function associated with an enumeration literal returns the corresponding value of the enumeration type.

NOTES

24 If an enumeration literal occurs in a context that does not otherwise suffice to determine the type of the literal, then qualification by the name of the enumeration type is one way to resolve the ambiguity (see 4.7).

Examples

Examples of enumeration types and subtypes:

```
type Day      is (Mon, Tue, Wed, Thu, Fri, Sat, Sun);
type Suit     is (Clubs, Diamonds, Hearts, Spades);
type Gender   is (M, F);
type Level    is (Low, Medium, Urgent);
type Color    is (White, Red, Yellow, Green, Blue, Brown, Black);
type Light    is (Red, Amber, Green); -- Red and Green are overloaded
type Hexa    is ('A', 'B', 'C', 'D', 'E', 'F');
type Mixed   is ('A', 'B', '*', B, None, '?', '%');
subtype Weekday is Day  range Mon .. Fri;
subtype Major   is Suit  range Hearts .. Spades;
subtype Rainbow is Color range Red .. Blue; -- the Color Red, not the Light
```

3.5.2 Character Types

Static Semantics

An enumeration type is said to be a *character type* if at least one of its enumeration literals is a `character_literal`.

The predefined type `Character` is a character type whose values correspond to the 256 code positions of Row 00 (also known as Latin-1) of the ISO/IEC 10646:2003 Basic Multilingual Plane (BMP). Each of the graphic characters of Row 00 of the BMP has a corresponding `character_literal` in `Character`. Each of the nongraphic positions of Row 00 (0000-001F and 007F-009F) has a corresponding language-defined name, which is not usable as an enumeration literal, but which is usable with the attributes `Image`, `Wide_Image`, `Wide_Wide_Image`, `Value`, `Wide_Value`, and `Wide_Wide_Value`; these names are given in the definition of type `Character` in A.1, “The Package Standard”, but are set in *italics*.

- 3/2 The predefined type `Wide_Character` is a character type whose values correspond to the 65536 code positions of the ISO/IEC 10646:2003 Basic Multilingual Plane (BMP). Each of the graphic characters of the BMP has a corresponding `character_literal` in `Wide_Character`. The first 256 values of `Wide_Character` have the same `character_literal` or language-defined name as defined for `Character`. Each of the `graphic_characters` has a corresponding `character_literal`.
- 3.1/2 The predefined type `Wide_Wide_Character` is a character type whose values correspond to the 2147483648 code positions of the ISO/IEC 10646:2003 character set. Each of the `graphic_characters` has a corresponding `character_literal` in `Wide_Wide_Character`. The first 65536 values of `Wide_Wide_Character` have the same `character_literal` or language-defined name as defined for `Wide_Character`.
- 3.2/2 The characters whose code position is larger than 16#FF# and which are not `graphic_characters` have language-defined names which are formed by appending to the string "Hex_" the representation of their code position in hexadecimal as eight extended digits. As with other language-defined names, these names are usable only with the attributes `(Wide_)Wide_Image` and `(Wide_)Wide_Value`; they are not usable as enumeration literals.

Implementation Permissions

- 4/2 *This paragraph was deleted.*

Implementation Advice

- 5/2 *This paragraph was deleted.*

NOTES

- 6 25 The language-defined library package `Characters.Latin_1` (see A.3.3) includes the declaration of constants denoting control characters, lower case characters, and special characters of the predefined type `Character`.
- 7 26 A conventional character set such as *EBCDIC* can be declared as a character type; the internal codes of the characters can be specified by an `enumeration_representation_clause` as explained in clause 13.4.

Examples

- 8 Example of a character type:

```
type Roman_Digit is ('I', 'V', 'X', 'L', 'C', 'D', 'M');
```

3.5.3 Boolean Types

Static Semantics

- 1 There is a predefined enumeration type named `Boolean`, declared in the visible part of package `Standard`. It has the two enumeration literals `False` and `True` ordered with the relation `False < True`. Any descendant of the predefined type `Boolean` is called a *boolean type*.

3.5.4 Integer Types

- 1 An `integer_type_definition` defines an integer type; it defines either a *signed* integer type, or a *modular* integer type. The base range of a signed integer type includes at least the values of the specified range. A modular type is an integer type with all arithmetic modulo a specified positive *modulus*; such a type corresponds to an unsigned type with wrap-around semantics.

Syntax

- 2 `integer_type_definition ::= signed_integer_type_definition | modular_type_definition`

signed_integer_type_definition ::= **range** static_simple_expression .. static_simple_expression
 modular_type_definition ::= **mod** static_expression

Name Resolution Rules

Each simple_expression in a signed_integer_type_definition is expected to be of any integer type; they need not be of the same type. The expression in a modular_type_definition is likewise expected to be of any integer type.

Legality Rules

The simple_expressions of a signed_integer_type_definition shall be static, and their values shall be in the range System.Min_Int .. System.Max_Int.

The expression of a modular_type_definition shall be static, and its value (the *modulus*) shall be positive, and shall be no greater than System.Max_Binary_Modulus if a power of 2, or no greater than System.Max_Nonbinary_Modulus if not.

Static Semantics

The set of values for a signed integer type is the (infinite) set of mathematical integers, though only values of the base range of the type are fully supported for run-time operations. The set of values for a modular integer type are the values from 0 to one less than the modulus, inclusive.

A signed_integer_type_definition defines an integer type whose base range includes at least the values of the simple_expressions and is symmetric about zero, excepting possibly an extra negative value. A signed_integer_type_definition also defines a constrained first subtype of the type, with a range whose bounds are given by the values of the simple_expressions, converted to the type being defined.

A modular_type_definition defines a modular type whose base range is from zero to one less than the given modulus. A modular_type_definition also defines a constrained first subtype of the type with a range that is the same as the base range of the type.

There is a predefined signed integer subtype named Integer, declared in the visible part of package Standard. It is constrained to the base range of its type.

Integer has two predefined subtypes, declared in the visible part of package Standard:

```
subtype Natural is Integer range 0 .. Integer'Last;
subtype Positive is Integer range 1 .. Integer'Last;
```

A type defined by an integer_type_definition is implicitly derived from root_integer, an anonymous predefined (specific) integer type, whose base range is System.Min_Int .. System.Max_Int. However, the base range of the new type is not inherited from root_integer, but is instead determined by the range or modulus specified by the integer_type_definition. Integer literals are all of the type universal_integer, the universal type (see 3.4.1) for the class rooted at root_integer, allowing their use with the operations of any integer type.

The position number of an integer value is equal to the value.

For every modular subtype S, the following attributes are defined:

S'Mod S'Mod denotes a function with the following specification:

```
function S'Mod (Arg : universal_integer)
               return S'Base
```

This function returns Arg mod S'Modulus, as a value of the type of S.

- 17 S'Modulus S'Modulus yields the modulus of the type of S, as a value of the type *universal_integer*.

Dynamic Semantics

- 18 The elaboration of an *integer_type_definition* creates the integer type and its first subtype.
- 19 For a modular type, if the result of the execution of a predefined operator (see 4.5) is outside the base range of the type, the result is reduced modulo the modulus of the type to a value that is within the base range of the type.
- 20 For a signed integer type, the exception *Constraint_Error* is raised by the execution of an operation that cannot deliver the correct result because it is outside the base range of the type. For any integer type, *Constraint_Error* is raised by the operators "/" , "rem" , and "mod" if the right operand is zero.

Implementation Requirements

- 21 In an implementation, the range of Integer shall include the range $-2^{**}15+1 \dots +2^{**}15-1$.
- 22 If Long_Integer is predefined for an implementation, then its range shall include the range $-2^{**}31+1 \dots +2^{**}31-1$.
- 23 System.Max_Binary_Modulus shall be at least $2^{**}16$.

Implementation Permissions

- 24 For the execution of a predefined operation of a signed integer type, the implementation need not raise *Constraint_Error* if the result is outside the base range of the type, so long as the correct result is produced.
- 25 An implementation may provide additional predefined signed integer types, declared in the visible part of Standard, whose first subtypes have names of the form Short_Integer, Long_Integer, Short_Short_Integer, Long_Long_Integer, etc. Different predefined integer types are allowed to have the same base range. However, the range of Integer should be no wider than that of Long_Integer. Similarly, the range of Short_Integer (if provided) should be no wider than Integer. Corresponding recommendations apply to any other predefined integer types. There need not be a named integer type corresponding to each distinct base range supported by an implementation. The range of each first subtype should be the base range of its type.
- 26 An implementation may provide *nonstandard integer types*, descendants of *root_integer* that are declared outside of the specification of package Standard, which need not have all the standard characteristics of a type defined by an *integer_type_definition*. For example, a nonstandard integer type might have an asymmetric base range or it might not be allowed as an array or loop index (a very long integer). Any type descended from a nonstandard integer type is also nonstandard. An implementation may place arbitrary restrictions on the use of such types; it is implementation defined whether operators that are predefined for "any integer type" are defined for a particular nonstandard integer type. In any case, such types are not permitted as *explicit_generic_actual_parameters* for formal scalar types — see 12.5.2.
- 27 For a one's complement machine, the high bound of the base range of a modular type whose modulus is one less than a power of 2 may be equal to the modulus, rather than one less than the modulus. It is implementation defined for which powers of 2, if any, this permission is exercised.
- 27.1/1 For a one's complement machine, implementations may support non-binary modulus values greater than System.Max_Nonbinary_Modulus. It is implementation defined which specific values greater than System.Max_Nonbinary_Modulus, if any, are supported.

Implementation Advice

An implementation should support Long_Integer in addition to Integer if the target machine supports 32-bit (or longer) arithmetic. No other named integer subtypes are recommended for package Standard. Instead, appropriate named integer subtypes should be provided in the library package Interfaces (see B.2).

An implementation for a two's complement machine should support modular types with a binary modulus up to System.Max_Int*2+2. An implementation should support a nonbinary modulus up to Integer'Last.

NOTES

27 Integer literals are of the anonymous predefined integer type *universal_integer*. Other integer types have no literals. However, the overload resolution rules (see 8.6, “The Context of Overload Resolution”) allow expressions of the type *universal_integer* whenever an integer type is expected.

28 The same arithmetic operators are predefined for all signed integer types defined by a *signed_integer_type_definition* (see 4.5, “Operators and Expression Evaluation”). For modular types, these same operators are predefined, plus bit-wise logical operators (**and**, **or**, **xor**, and **not**). In addition, for the unsigned types declared in the language-defined package Interfaces (see B.2), functions are defined that provide bit-wise shifting and rotating.

29 Modular types match a *generic_formal_parameter_declaration* of the form "type T is mod <>;" signed integer types match "type T is range <>;" (see 12.5.2).

*Examples**Examples of integer types and subtypes:*

```
type Page_Num is range 1 .. 2_000;
type Line_Size is range 1 .. Max_Line_Size;
subtype Small_Int is Integer range -10 .. 10;
subtype Column_Ptr is Line_Size range 1 .. 10;
subtype Buffer_Size is Integer range 0 .. Max;
type Byte is mod 256; -- an unsigned byte
type Hash_Index is mod 97; -- modulus is prime
```

3.5.5 Operations of Discrete Types*Static Semantics*

For every discrete subtype S, the following attributes are defined:

S'Pos S'Pos denotes a function with the following specification:

```
function S'Pos(Arg : S'Base)
    return universal_integer
```

This function returns the position number of the value of Arg, as a value of type *universal_integer*.

S'Val S'Val denotes a function with the following specification:

```
function S'Val(Arg : universal_integer)
    return S'Base
```

This function returns a value of the type of S whose position number equals the value of Arg. For the evaluation of a call on S'Val, if there is no value in the base range of its type with the given position number, Constraint_Error is raised.

Implementation Advice

For the evaluation of a call on S'Pos for an enumeration subtype, if the value of the operand does not correspond to the internal code for any enumeration literal of its type (perhaps due to an uninitialized variable), then the implementation should raise Program_Error. This is particularly important for enumeration types with noncontiguous internal codes specified by an *enumeration_representation_clause*.

NOTES

9 30 Indexing and loop iteration use values of discrete types.

10 31 The predefined operations of a discrete type include the assignment operation, qualification, the membership tests, and the relational operators; for a boolean type they include the short-circuit control forms and the logical operators; for an integer type they include type conversion to and from other numeric types, as well as the binary and unary adding operators – and +, the multiplying operators, the unary operator **abs**, and the exponentiation operator. The assignment operation is described in 5.2. The other predefined operations are described in Section 4.

11 32 As for all types, objects of a discrete type have **Size** and **Address** attributes (see 13.3).

12 33 For a subtype of a discrete type, the result delivered by the attribute **Val** might not belong to the subtype; similarly, the actual parameter of the attribute **Pos** need not belong to the subtype. The following relations are satisfied (in the absence of an exception) by these attributes:

13 $S'Val(S'Pos(X)) = X$
 $S'Pos(S'Val(N)) = N$

*Examples**Examples of attributes of discrete subtypes:*

15 -- For the types and subtypes declared in subclause 3.5.1 the following hold:

16 -- Color'First = White, Color'Last = Black
-- Rainbow'First = Red, Rainbow'Last = Blue
17 -- Color'Succ(Blue) = Rainbow'Succ(Blue) = Brown
-- Color'Pos(Blue) = Rainbow'Pos(Blue) = 4
-- Color'Val(0) = Rainbow'Val(0) = White

3.5.6 Real Types

1 Real types provide approximations to the real numbers, with relative bounds on errors for floating point types, and with absolute bounds for fixed point types.

Syntax

2 **real_type_definition ::=**
floating_point_definition | fixed_point_definition

Static Semantics

3 A type defined by a **real_type_definition** is implicitly derived from **root_real**, an anonymous predefined (specific) real type. Hence, all real types, whether floating point or fixed point, are in the derivation class rooted at **root_real**.

4 Real literals are all of the type **universal_real**, the universal type (see 3.4.1) for the class rooted at **root_real**, allowing their use with the operations of any real type. Certain multiplying operators have a result type of **universal_fixed** (see 4.5.5), the universal type for the class of fixed point types, allowing the result of the multiplication or division to be used where any specific fixed point type is expected.

Dynamic Semantics

5 The elaboration of a **real_type_definition** consists of the elaboration of the **floating_point_definition** or the **fixed_point_definition**.

Implementation Requirements

6 An implementation shall perform the run-time evaluation of a use of a predefined operator of **root_real** with an accuracy at least as great as that of any floating point type definable by a **floating_point_definition**.

Implementation Permissions

For the execution of a predefined operation of a real type, the implementation need not raise Constraint_Error if the result is outside the base range of the type, so long as the correct result is produced, or the Machine_Overflows attribute of the type is False (see G.2).

An implementation may provide *nonstandard real types*, descendants of *root_real* that are declared outside of the specification of package Standard, which need not have all the standard characteristics of a type defined by a *real_type_definition*. For example, a nonstandard real type might have an asymmetric or unsigned base range, or its predefined operations might wrap around or “saturate” rather than overflow (modular or saturating arithmetic), or it might not conform to the accuracy model (see G.2). Any type descended from a nonstandard real type is also nonstandard. An implementation may place arbitrary restrictions on the use of such types; it is implementation defined whether operators that are predefined for “any real type” are defined for a particular nonstandard real type. In any case, such types are not permitted as *explicit_generic_actual_parameters* for formal scalar types — see 12.5.2.

NOTES

34 As stated, real literals are of the anonymous predefined real type *universal_real*. Other real types have no literals. However, the overload resolution rules (see 8.6) allow expressions of the type *universal_real* whenever a real type is expected.

3.5.7 Floating Point Types

For floating point types, the error bound is specified as a relative precision by giving the required minimum number of significant decimal digits.

Syntax

```

floating_point_definition ::=           2
  digits static_expression [real_range_specification]
real_range_specification ::=           3
  range static_simple_expression .. static_simple_expression

```

Name Resolution Rules

The *requested decimal precision*, which is the minimum number of significant decimal digits required for the floating point type, is specified by the value of the expression given after the reserved word **digits**. This expression is expected to be of any integer type.

Each *simple_expression* of a *real_range_specification* is expected to be of any real type; the types need not be the same.

Legality Rules

The requested decimal precision shall be specified by a static expression whose value is positive and no greater than System.Max_Base_Digits. Each *simple_expression* of a *real_range_specification* shall also be static. If the *real_range_specification* is omitted, the requested decimal precision shall be no greater than System.Max_Digits.

A *floating_point_definition* is illegal if the implementation does not support a floating point type that satisfies the requested decimal precision and range.

Static Semantics

The set of values for a floating point type is the (infinite) set of rational numbers. The *machine numbers* of a floating point type are the values of the type that can be represented exactly in every unconstrained

variable of the type. The base range (see 3.5) of a floating point type is symmetric around zero, except that it can include some extra negative values in some implementations.

- 9 The *base decimal precision* of a floating point type is the number of decimal digits of precision representable in objects of the type. The *safe range* of a floating point type is that part of its base range for which the accuracy corresponding to the base decimal precision is preserved by all predefined operations.
- 10 A *floating_point_definition* defines a floating point type whose base decimal precision is no less than the requested decimal precision. If a *real_range_specification* is given, the safe range of the floating point type (and hence, also its base range) includes at least the values of the simple expressions given in the *real_range_specification*. If a *real_range_specification* is not given, the safe (and base) range of the type includes at least the values of the range $-10.0^{**}(4*D) .. +10.0^{**}(4*D)$ where D is the requested decimal precision. The safe range might include other values as well. The attributes *Safe_First* and *Safe_Last* give the actual bounds of the safe range.
- 11 A *floating_point_definition* also defines a first subtype of the type. If a *real_range_specification* is given, then the subtype is constrained to a range whose bounds are given by a conversion of the values of the *simple_expressions* of the *real_range_specification* to the type being defined. Otherwise, the subtype is unconstrained.
- 12 There is a predefined, unconstrained, floating point subtype named *Float*, declared in the visible part of package Standard.

Dynamic Semantics

- 13 The elaboration of a *floating_point_definition* creates the floating point type and its first subtype.

Implementation Requirements

- 14 In an implementation that supports floating point types with 6 or more digits of precision, the requested decimal precision for *Float* shall be at least 6.
- 15 If *Long_Float* is predefined for an implementation, then its requested decimal precision shall be at least 11.

Implementation Permissions

- 16 An implementation is allowed to provide additional predefined floating point types, declared in the visible part of Standard, whose (unconstrained) first subtypes have names of the form *Short_Float*, *Long_Float*, *Short_Short_Float*, *Long_Long_Float*, etc. Different predefined floating point types are allowed to have the same base decimal precision. However, the precision of *Float* should be no greater than that of *Long_Float*. Similarly, the precision of *Short_Float* (if provided) should be no greater than *Float*. Corresponding recommendations apply to any other predefined floating point types. There need not be a named floating point type corresponding to each distinct base decimal precision supported by an implementation.

Implementation Advice

- 17 An implementation should support *Long_Float* in addition to *Float* if the target machine supports 11 or more digits of precision. No other named floating point subtypes are recommended for package Standard. Instead, appropriate named floating point subtypes should be provided in the library package Interfaces (see B.2).

NOTES

- 18 35 If a floating point subtype is unconstrained, then assignments to variables of the subtype involve only *Overflow_Checks*, never *Range_Checks*.

*Examples**Examples of floating point types and subtypes:*

```

type Coefficient is digits 10 range -1.0 .. 1.0;
type Real is digits 8;
type Mass is digits 7 range 0.0 .. 1.0E35;
subtype Probability is Real range 0.0 .. 1.0;    -- a subtype with a smaller range

```

3.5.8 Operations of Floating Point Types

Static Semantics

The following attribute is defined for every floating point subtype S:

S'Digits S'Digits denotes the requested decimal precision for the subtype S. The value of this attribute is of the type *universal_integer*. The requested decimal precision of the base subtype of a floating point type T is defined to be the largest value of d for which
 $\text{ceiling}(d * \log(10) / \log(T'\text{Machine_Radix})) + g \leq T'\text{Model_Mantissa}$
where g is 0 if Machine_Radix is a positive power of 10 and 1 otherwise.

NOTES

36 The predefined operations of a floating point type include the assignment operation, qualification, the membership tests, and explicit conversion to and from other numeric types. They also include the relational operators and the following predefined arithmetic operators: the binary and unary adding operators – and +, certain multiplying operators, the unary operator **abs**, and the exponentiation operator.

37 As for all types, objects of a floating point type have Size and Address attributes (see 13.3). Other attributes of floating point types are defined in A.5.3.

3.5.9 Fixed Point Types

A fixed point type is either an ordinary fixed point type, or a decimal fixed point type. The error bound of a fixed point type is specified as an absolute value, called the *delta* of the fixed point type.*Syntax*

```

fixed_point_definition ::= ordinary_fixed_point_definition | decimal_fixed_point_definition
ordinary_fixed_point_definition ::= delta static_expression real_range_specification
decimal_fixed_point_definition ::= delta static_expression digits static_expression [real_range_specification]
digits_constraint ::= digits static_expression [range_constraint]

```

*Name Resolution Rules*For a type defined by a *fixed_point_definition*, the *delta* of the type is specified by the value of the expression given after the reserved word **delta**; this expression is expected to be of any real type. For a type defined by a *decimal_fixed_point_definition* (a *decimal* fixed point type), the number of significant decimal digits for its first subtype (the *digits* of the first subtype) is specified by the expression given after the reserved word **digits**; this expression is expected to be of any integer type.*Legality Rules*In a *fixed_point_definition* or *digits_constraint*, the expressions given after the reserved words **delta** and **digits** shall be static; their values shall be positive.

- 8/2 The set of values of a fixed point type comprise the integral multiples of a number called the *small* of the type. The *machine numbers* of a fixed point type are the values of the type that can be represented exactly in every unconstrained variable of the type. For a type defined by an `ordinary_fixed_point_definition` (an *ordinary* fixed point type), the *small* may be specified by an `attribute_definition_clause` (see 13.3); if so specified, it shall be no greater than the *delta* of the type. If not specified, the *small* of an ordinary fixed point type is an implementation-defined power of two less than or equal to the *delta*.
- 9 For a decimal fixed point type, the *small* equals the *delta*; the *delta* shall be a power of 10. If a `real_range_specification` is given, both bounds of the range shall be in the range $-(10^{**}d\text{igits}-1)*\delta\text{elta} .. +(10^{**}d\text{igits}-1)*\delta\text{elta}$.
- 10 A `fixed_point_definition` is illegal if the implementation does not support a fixed point type with the given *small* and specified range or *digits*.
- 11 For a `subtype_indication` with a `digits_constraint`, the `subtype_mark` shall denote a decimal fixed point subtype.

Static Semantics

- 12 The base range (see 3.5) of a fixed point type is symmetric around zero, except possibly for an extra negative value in some implementations.
- 13 An `ordinary_fixed_point_definition` defines an ordinary fixed point type whose base range includes at least all multiples of *small* that are between the bounds specified in the `real_range_specification`. The base range of the type does not necessarily include the specified bounds themselves. An `ordinary_fixed_point_definition` also defines a constrained first subtype of the type, with each bound of its range given by the closer to zero of:
- the value of the conversion to the fixed point type of the corresponding expression of the `real_range_specification`;
 - the corresponding bound of the base range.
- 14 A `decimal_fixed_point_definition` defines a decimal fixed point type whose base range includes at least the range $-(10^{**}d\text{igits}-1)*\delta\text{elta} .. +(10^{**}d\text{igits}-1)*\delta\text{elta}$. A `decimal_fixed_point_definition` also defines a constrained first subtype of the type. If a `real_range_specification` is given, the bounds of the first subtype are given by a conversion of the values of the expressions of the `real_range_specification`. Otherwise, the range of the first subtype is $-(10^{**}d\text{igits}-1)*\delta\text{elta} .. +(10^{**}d\text{igits}-1)*\delta\text{elta}$.

Dynamic Semantics

- 17 The elaboration of a `fixed_point_definition` creates the fixed point type and its first subtype.
- 18 For a `digits_constraint` on a decimal fixed point subtype with a given *delta*, if it does not have a `range_constraint`, then it specifies an implicit range $-(10^{**}D-1)*\delta\text{elta} .. +(10^{**}D-1)*\delta\text{elta}$, where *D* is the value of the expression. A `digits_constraint` is *compatible* with a decimal fixed point subtype if the value of the expression is no greater than the *digits* of the subtype, and if it specifies (explicitly or implicitly) a range that is compatible with the subtype.
- 19 The elaboration of a `digits_constraint` consists of the elaboration of the `range_constraint`, if any. If a `range_constraint` is given, a check is made that the bounds of the range are both in the range $-(10^{**}D-1)*\delta\text{elta} .. +(10^{**}D-1)*\delta\text{elta}$, where *D* is the value of the (static) expression given after the reserved word **digits**. If this check fails, `Constraint_Error` is raised.

Implementation Requirements

The implementation shall support at least 24 bits of precision (including the sign bit) for fixed point types. 20

Implementation Permissions

Implementations are permitted to support only *smalls* that are a power of two. In particular, all decimal fixed point type declarations can be disallowed. Note however that conformance with the Information Systems Annex requires support for decimal *smalls*, and decimal fixed point type declarations with *digits* up to at least 18. 21

NOTES

38 The base range of an ordinary fixed point type need not include the specified bounds themselves so that the range specification can be given in a natural way, such as: 22

```
type Fraction is delta 2.0**(-15) range -1.0 .. 1.0; 23
```

With 2's complement hardware, such a type could have a signed 16-bit representation, using 1 bit for the sign and 15 bits for fraction, resulting in a base range of -1.0 .. 1.0-2.0**(-15). 24

*Examples**Examples of fixed point types and subtypes:* 25

```
type Volt is delta 0.125 range 0.0 .. 255.0;
  -- A pure fraction which requires all the available
  -- space in a word can be declared as the type Fraction:
type Fraction is delta System.Fine_Delta range -1.0 .. 1.0;
  -- Fraction'Last = 1.0 - System.Fine_Delta
type Money is delta 0.01 digits 15;  -- decimal fixed point
subtype Salary is Money digits 10;
  -- Money'Last = 10.0**13 - 0.01, Salary'Last = 10.0**8 - 0.01 28
```

3.5.10 Operations of Fixed Point Types

Static Semantics

The following attributes are defined for every fixed point subtype S: 1

S'Small S'Small denotes the *small* of the type of S. The value of this attribute is of the type *universal_real*. Small may be specified for nonderived ordinary fixed point types via an *attribute_definition_clause* (see 13.3); the expression of such a clause shall be static. 2/1

S'Delta S'Delta denotes the *delta* of the fixed point subtype S. The value of this attribute is of the type *universal_real*. 3

S'Fore S'Fore yields the minimum number of characters needed before the decimal point for the decimal representation of any value of the subtype S, assuming that the representation does not include an exponent, but includes a one-character prefix that is either a minus sign or a space. (This minimum number does not include superfluous zeros or underlines, and is at least 2.) The value of this attribute is of the type *universal_integer*. 4

S'Aft S'Aft yields the number of decimal digits needed after the decimal point to accommodate the *delta* of the subtype S, unless the *delta* of the subtype S is greater than 0.1, in which case the attribute yields the value one. (S'Aft is the smallest positive integer N for which $(10^{**}N)*S'Delta$ is greater than or equal to one.) The value of this attribute is of the type *universal_integer*. 5

The following additional attributes are defined for every decimal fixed point subtype S: 6

- 7 S'Digits S'Digits denotes the *digits* of the decimal fixed point subtype S, which corresponds to the number of decimal digits that are representable in objects of the subtype. The value of this attribute is of the type *universal_integer*. Its value is determined as follows:
- 8 • For a first subtype or a subtype defined by a *subtype_indication* with a *digits_constraint*, the digits is the value of the expression given after the reserved word **digits**;
- 9 • For a subtype defined by a *subtype_indication* without a *digits_constraint*, the digits of the subtype is the same as that of the subtype denoted by the *subtype_mark* in the *subtype_indication*.
- 10 • The digits of a base subtype is the largest integer D such that the range $-(10^{**D-1})*\delta .. +(10^{**D-1})*\delta$ is included in the base range of the type.
- 11 S'Scale S'Scale denotes the *scale* of the subtype S, defined as the value N such that $S'Delta = 10.0^{**(-N)}$. The scale indicates the position of the point relative to the rightmost significant digits of values of subtype S. The value of this attribute is of the type *universal_integer*.
- 12 S'Round S'Round denotes a function with the following specification:
- ```
13 function S'Round(X : universal_real)
14 return S'Base
```
- 14        The function returns the value obtained by rounding X (away from 0, if X is midway between two values of the type of S).
- 15        NOTES  
 39 All subtypes of a fixed point type will have the same value for the Delta attribute, in the absence of *delta\_constraints* (see J.3).
- 16        40 S'Scale is not always the same as S'Aft for a decimal subtype; for example, if S'Delta = 1.0 then S'Aft is 1 while S'Scale is 0.
- 17        41 The predefined operations of a fixed point type include the assignment operation, qualification, the membership tests, and explicit conversion to and from other numeric types. They also include the relational operators and the following predefined arithmetic operators: the binary and unary adding operators – and +, multiplying operators, and the unary operator **abs**.
- 18        42 As for all types, objects of a fixed point type have Size and Address attributes (see 13.3). Other attributes of fixed point types are defined in A.5.4.

## 3.6 Array Types

- 1    An *array* object is a composite object consisting of components which all have the same subtype. The name for a component of an array uses one or more index values belonging to specified discrete types. The value of an array object is a composite value consisting of the values of the components.

### Syntax

```
2 array_type_definition ::=
3 unconstrained_array_definition | constrained_array_definition
4
3 unconstrained_array_definition ::=
5 array(index_subtype_definition {, index_subtype_definition}) of component_definition
6
4 index_subtype_definition ::= subtype_mark range <>
5
5 constrained_array_definition ::=
6 array (discrete_subtype_definition {, discrete_subtype_definition}) of component_definition
6
6 discrete_subtype_definition ::= discrete_subtype_indication | range
```

component\_definition ::=  
 [aliased] subtype\_indication  
 | [aliased] access\_definition

*Name Resolution Rules*

For a discrete\_subtype\_definition that is a range, the range shall resolve to be of some specific discrete type; which discrete type shall be determined without using any context other than the bounds of the range itself (plus the preference for root\_integer — see 8.6).

*Legality Rules*

Each index\_subtype\_definition or discrete\_subtype\_definition in an array\_type\_definition defines an index subtype; its type (the *index type*) shall be discrete.

The subtype defined by the subtype\_indication of a component\_definition (the *component subtype*) shall be a definite subtype.

*This paragraph was deleted.*

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*Static Semantics*

An array is characterized by the number of indices (the *dimensionality* of the array), the type and position of each index, the lower and upper bounds for each index, and the subtype of the components. The order of the indices is significant.

A one-dimensional array has a distinct component for each possible index value. A multidimensional array has a distinct component for each possible sequence of index values that can be formed by selecting one value for each index position (in the given order). The possible values for a given index are all the values between the lower and upper bounds, inclusive; this range of values is called the *index range*. The *bounds* of an array are the bounds of its index ranges. The *length* of a dimension of an array is the number of values of the index range of the dimension (zero for a null range). The *length* of a one-dimensional array is the length of its only dimension.

An array\_type\_definition defines an array type and its first subtype. For each object of this array type, the number of indices, the type and position of each index, and the subtype of the components are as in the type definition; the values of the lower and upper bounds for each index belong to the corresponding index subtype of its type, except for null arrays (see 3.6.1).

An unconstrained\_array\_definition defines an array type with an unconstrained first subtype. Each index\_subtype\_definition defines the corresponding index subtype to be the subtype denoted by the subtype\_mark. The compound delimiter <> (called a *box*) of an index\_subtype\_definition stands for an undefined range (different objects of the type need not have the same bounds).

A constrained\_array\_definition defines an array type with a constrained first subtype. Each discrete\_subtype\_definition defines the corresponding index subtype, as well as the corresponding index range for the constrained first subtype. The *constraint* of the first subtype consists of the bounds of the index ranges.

The discrete subtype defined by a discrete\_subtype\_definition is either that defined by the subtype\_indication, or a subtype determined by the range as follows:

- If the type of the range resolves to root\_integer, then the discrete\_subtype\_definition defines a subtype of the predefined type Integer with bounds given by a conversion to Integer of the bounds of the range;
- Otherwise, the discrete\_subtype\_definition defines a subtype of the type of the range, with the bounds given by the range.

- 20 The `component_definition` of an `array_type_definition` defines the nominal subtype of the components. If the reserved word `aliased` appears in the `component_definition`, then each component of the array is aliased (see 3.10).

*Dynamic Semantics*

- 21 The elaboration of an `array_type_definition` creates the array type and its first subtype, and consists of the elaboration of any `discrete_subtype_definitions` and the `component_definition`.
- 22/2 The elaboration of a `discrete_subtype_definition` that does not contain any per-object expressions creates the discrete subtype, and consists of the elaboration of the `subtype_indication` or the evaluation of the range. The elaboration of a `discrete_subtype_definition` that contains one or more per-object expressions is defined in 3.8. The elaboration of a `component_definition` in an `array_type_definition` consists of the elaboration of the `subtype_indication` or `access_definition`. The elaboration of any `discrete_subtype_definitions` and the elaboration of the `component_definition` are performed in an arbitrary order.

**NOTES**

- 23 43 All components of an array have the same subtype. In particular, for an array of components that are one-dimensional arrays, this means that all components have the same bounds and hence the same length.
- 24 44 Each elaboration of an `array_type_definition` creates a distinct array type. A consequence of this is that each object whose `object_declaration` contains an `array_type_definition` is of its own unique type.

*Examples*

- 25 Examples of type declarations with unconstrained array definitions:

```
26 type Vector is array(Integer range <>) of Real;
 type Matrix is array(Integer range <>, Integer range <>) of Real;
 type Bit_Vector is array(Integer range <>) of Boolean;
 type Roman is array(Positive range <>) of Roman_Digit; -- see 3.5.2
```

- 27 Examples of type declarations with constrained array definitions:

```
28 type Table is array(1 .. 10) of Integer;
 type Schedule is array(Day) of Boolean;
 type Line is array(1 .. Max_Line_Size) of Character;
```

- 29 Examples of object declarations with array type definitions:

```
30/2 Grid : array(1 .. 80, 1 .. 100) of Boolean;
 Mix : array(Color range Red .. Green) of Boolean;
 Msg_Table : constant array(Error_Code) of access constant String :=
 (Too_Big => new String("Result too big"), Too_Small => ...);
 Page : array(Positive range <>) of Line := -- an array of arrays
 (1 | 50 => Line'(1 | Line'Last => '+', others => '-'), -- see 4.3.3
 2 .. 49 => Line'(1 | Line'Last => '|', others => ' '));
-- Page is constrained by its initial value to (1..50)
```

### 3.6.1 Index Constraints and Discrete Ranges

- 1 An `index_constraint` determines the range of possible values for every index of an array subtype, and thereby the corresponding array bounds.

*Syntax*

- 2 `index_constraint ::= (discrete_range {, discrete_range})`
- 3 `discrete_range ::= discrete_subtype_indication | range`

*Name Resolution Rules*

The type of a `discrete_range` is the type of the subtype defined by the `subtype_indication`, or the type of the `range`. For an `index_constraint`, each `discrete_range` shall resolve to be of the type of the corresponding index.

*Legality Rules*

An `index_constraint` shall appear only in a `subtype_indication` whose `subtype_mark` denotes either an unconstrained array subtype, or an unconstrained access subtype whose designated subtype is an unconstrained array subtype; in either case, the `index_constraint` shall provide a `discrete_range` for each index of the array type.

*Static Semantics*

A `discrete_range` defines a range whose bounds are given by the `range`, or by the range of the subtype defined by the `subtype_indication`.

*Dynamic Semantics*

An `index_constraint` is *compatible* with an unconstrained array subtype if and only if the index range defined by each `discrete_range` is compatible (see 3.5) with the corresponding index subtype. If any of the `discrete_ranges` defines a null range, any array thus constrained is a *null array*, having no components. An array value *satisfies* an `index_constraint` if at each index position the array value and the `index_constraint` have the same index bounds.

The elaboration of an `index_constraint` consists of the evaluation of the `discrete_range`(s), in an arbitrary order. The evaluation of a `discrete_range` consists of the elaboration of the `subtype_indication` or the evaluation of the `range`.

## NOTES

45 The elaboration of a `subtype_indication` consisting of a `subtype_mark` followed by an `index_constraint` checks the compatibility of the `index_constraint` with the `subtype_mark` (see 3.2.2).

46 Even if an array value does not satisfy the index constraint of an array subtype, `Constraint_Error` is not raised on conversion to the array subtype, so long as the length of each dimension of the array value and the array subtype match. See 4.6.

*Examples*

*Examples of array declarations including an index constraint:*

```
Board : Matrix(1 .. 8, 1 .. 8); -- see 3.6
Rectangle : Matrix(1 .. 20, 1 .. 30);
Inverse : Matrix(1 .. N, 1 .. N); -- N need not be static
Filter : Bit_Vector(0 .. 31);
```

*Example of array declaration with a constrained array subtype:*

```
My_Schedule : Schedule; -- all arrays of type Schedule have the same bounds
```

*Example of record type with a component that is an array:*

```
type Var_Line(Length : Natural) is
 record
 Image : String(1 .. Length);
 end record;
Null_Line : Var_Line(0); -- Null_Line.Image is a null array
```

## 3.6.2 Operations of Array Types

### *Legality Rules*

- 1 The argument N used in the `attribute_designators` for the N-th dimension of an array shall be a static expression of some integer type. The value of N shall be positive (nonzero) and no greater than the dimensionality of the array.

### *Static Semantics*

- 2/1 The following attributes are defined for a prefix A that is of an array type (after any implicit dereference), or denotes a constrained array subtype:

- 3 A'First      A'First denotes the lower bound of the first index range; its type is the corresponding index type.
- 4 A'First(N)    A'First(N) denotes the lower bound of the N-th index range; its type is the corresponding index type.
- 5 A'Last        A'Last denotes the upper bound of the first index range; its type is the corresponding index type.
- 6 A'Last(N)     A'Last(N) denotes the upper bound of the N-th index range; its type is the corresponding index type.
- 7 A'Range       A'Range is equivalent to the range A'First .. A'Last, except that the prefix A is only evaluated once.
- 8 A'Range(N)    A'Range(N) is equivalent to the range A'First(N) .. A'Last(N), except that the prefix A is only evaluated once.
- 9 A'Length      A'Length denotes the number of values of the first index range (zero for a null range); its type is *universal\_integer*.
- 10 A'Length(N)   A'Length(N) denotes the number of values of the N-th index range (zero for a null range); its type is *universal\_integer*.

### *Implementation Advice*

- 11 An implementation should normally represent multidimensional arrays in row-major order, consistent with the notation used for multidimensional array aggregates (see 4.3.3). However, if a **pragma** Convention(Fortran, ...) applies to a multidimensional array type, then column-major order should be used instead (see B.5, “Interfacing with Fortran”).

#### NOTES

12 47 The `attribute_references` A'First and A'First(1) denote the same value. A similar relation exists for the `attribute_references` A'Last, A'Range, and A'Length. The following relation is satisfied (except for a null array) by the above attributes if the index type is an integer type:

$$A'Length(N) = A'Last(N) - A'First(N) + 1$$

14 48 An array type is limited if its component type is limited (see 7.5).

15 49 The predefined operations of an array type include the membership tests, qualification, and explicit conversion. If the array type is not limited, they also include assignment and the predefined equality operators. For a one-dimensional array type, they include the predefined concatenation operators (if nonlimited) and, if the component type is discrete, the predefined relational operators; if the component type is boolean, the predefined logical operators are also included.

16/2 50 A component of an array can be named with an `indexed_component`. A value of an array type can be specified with an `array_aggregate`. For a one-dimensional array type, a slice of the array can be named; also, string literals are defined if the component type is a character type.

*Examples**Examples (using arrays declared in the examples of subclause 3.6.1):*

```
-- Filter'First = 0 Filter'Last = 31 Filter'Length = 32 17
-- Rectangle'Last(1) = 20 Rectangle'Last(2) = 30
```

### 3.6.3 String Types

*Static Semantics*A one-dimensional array type whose component type is a character type is called a *string* type.There are three predefined string types, *String*, *Wide\_String*, and *Wide\_Wide\_String*, each indexed by values of the predefined subtype *Positive*; these are declared in the visible part of package *Standard*:

```
subtype Positive is Integer range 1 .. Integer'Last; 3
type String is array(Positive range <>) of Character;
type Wide_String is array(Positive range <>) of Wide_Character;
type Wide_Wide_String is array(Positive range <>) of Wide_Wide_Character;
```

## NOTES

51 String literals (see 2.6 and 4.2) are defined for all string types. The concatenation operator `&` is predefined for string types, as for all nonlimited one-dimensional array types. The ordering operators `<`, `<=`, `>`, and `>=` are predefined for string types, as for all one-dimensional discrete array types; these ordering operators correspond to lexicographic order (see 4.5.2).*Examples**Examples of string objects:*

```
Stars : String(1 .. 120) := (1 .. 120 => '*');
Question : constant String := "How many characters?";
 -- Question'First = 1, Question'Last = 20
 -- Question'Length = 20 (the number of
characters)

Ask_Twice : String := Question & Question; -- constrained to (1..40)
Ninety_Six: constant Roman := "XCVI"; -- see 3.5.2 and 3.6
```

### 3.7 Discriminants

A composite type (other than an array or interface type) can have discriminants, which parameterize the type. A *known\_discriminant\_part* specifies the discriminants of a composite type. A discriminant of an object is a component of the object, and is either of a discrete type or an access type. An *unknown\_discriminant\_part* in the declaration of a view of a type specifies that the discriminants of the type are unknown for the given view; all subtypes of such a view are indefinite subtypes.*Syntax*

```
discriminant_part ::= unknown_discriminant_part | known_discriminant_part 2/2
unknown_discriminant_part ::= (<>)
known_discriminant_part ::= (discriminant_specification {; discriminant_specification})
discriminant_specification ::= defining_identifier_list : [null_exclusion] subtype_mark [:default_expression]
 | defining_identifier_list : access_definition [:default_expression]
default_expression ::= expression
```

*Name Resolution Rules*

- 7 The expected type for the `default_expression` of a `discriminant_specification` is that of the corresponding discriminant.

*Legality Rules*

- 8/2 A `discriminant_part` is only permitted in a declaration for a composite type that is not an array or interface type (this includes generic formal types). A type declared with a `known_discriminant_part` is called a *discriminated type*, as is a type that inherits (known) discriminants.
- 9/2 The subtype of a discriminant may be defined by an optional `null_exclusion` and a `subtype_mark`, in which case the `subtype_mark` shall denote a discrete or access subtype, or it may be defined by an `access_definition`. A discriminant that is defined by an `access_definition` is called an *access discriminant* and is of an anonymous access type.
- 9.1/2 `Default_expressions` shall be provided either for all or for none of the discriminants of a `known_discriminant_part`. No `default_expressions` are permitted in a `known_discriminant_part` in a declaration of a tagged type or a generic formal type.
- 10/2 A `discriminant_specification` for an access discriminant may have a `default_expression` only in the declaration for a task or protected type, or for a type that is a descendant of an explicitly limited record type. In addition to the places where Legality Rules normally apply (see 12.3), this rule applies also in the private part of an instance of a generic unit.
- 11/2 *This paragraph was deleted.*
- 12 For a type defined by a `derived_type_definition`, if a `known_discriminant_part` is provided in its declaration, then:
- 13 • The parent subtype shall be constrained;
  - 14 • If the parent type is not a tagged type, then each discriminant of the derived type shall be used in the constraint defining the parent subtype;
  - 15 • If a discriminant is used in the constraint defining the parent subtype, the subtype of the discriminant shall be statically compatible (see 4.9.1) with the subtype of the corresponding parent discriminant.
- 16 The type of the `default_expression`, if any, for an access discriminant shall be convertible to the anonymous access type of the discriminant (see 4.6).

*Static Semantics*

- 17 A `discriminant_specification` declares a discriminant; the `subtype_mark` denotes its subtype unless it is an access discriminant, in which case the discriminant's subtype is the anonymous access-to-variable subtype defined by the `access_definition`.
- 18 For a type defined by a `derived_type_definition`, each discriminant of the parent type is either inherited, constrained to equal some new discriminant of the derived type, or constrained to the value of an expression. When inherited or constrained to equal some new discriminant, the parent discriminant and the discriminant of the derived type are said to *correspond*. Two discriminants also correspond if there is some common discriminant to which they both correspond. A discriminant corresponds to itself as well. If a discriminant of a parent type is constrained to a specific value by a `derived_type_definition`, then that discriminant is said to be *specified* by that `derived_type_definition`.

A constraint that appears within the definition of a discriminated type *depends on a discriminant* of the type if it names the discriminant as a bound or discriminant value. A `component_definition` depends on a discriminant if its constraint depends on the discriminant, or on a discriminant that corresponds to it.

A component *depends on a discriminant* if:

- Its `component_definition` depends on the discriminant; or
- It is declared in a `variant_part` that is governed by the discriminant; or
- It is a component inherited as part of a `derived_type_definition`, and the constraint of the `parent_subtype_indication` depends on the discriminant; or
- It is a subcomponent of a component that depends on the discriminant.

Each value of a discriminated type includes a value for each component of the type that does not depend on a discriminant; this includes the discriminants themselves. The values of discriminants determine which other component values are present in the value of the discriminated type.

A type declared with a `known_discriminant_part` is said to have *known discriminants*; its first subtype is unconstrained. A type declared with an `unknown_discriminant_part` is said to have *unknown discriminants*. A type declared without a `discriminant_part` has no discriminants, unless it is a derived type; if derived, such a type has the same sort of discriminants (known, unknown, or none) as its parent (or ancestor) type. A tagged class-wide type also has unknown discriminants. Any subtype of a type with unknown discriminants is an unconstrained and indefinite subtype (see 3.2 and 3.3).

#### *Dynamic Semantics*

For an access discriminant, its `access_definition` is elaborated when the value of the access discriminant is defined: by evaluation of its `default_expression`, by elaboration of a `discriminant_constraint`, or by an assignment that initializes the enclosing object.

#### NOTES

52 If a discriminated type has `default_expressions` for its discriminants, then unconstrained variables of the type are permitted, and the values of the discriminants can be changed by an assignment to such a variable. If defaults are not provided for the discriminants, then all variables of the type are constrained, either by explicit constraint or by their initial value; the values of the discriminants of such a variable cannot be changed after initialization.

53 The `default_expression` for a discriminant of a type is evaluated when an object of an unconstrained subtype of the type is created.

54 Assignment to a discriminant of an object (after its initialization) is not allowed, since the name of a discriminant is a constant; neither `assignment_statements` nor assignments inherent in passing as an `in out` or `out` parameter are allowed. Note however that the value of a discriminant can be changed by assigning to the enclosing object, presuming it is an unconstrained variable.

55 A discriminant that is of a named access type is not called an access discriminant; that term is used only for discriminants defined by an `access_definition`.

#### *Examples*

*Examples of discriminated types:*

```
type Buffer(Size : Buffer_Size := 100) is
 record
 Pos : Buffer_Size := 0;
 Value : String(1 .. Size);
 end record;

type Matrix_Rec(Rows, Columns : Integer) is
 record
 Mat : Matrix(1 .. Rows, 1 .. Columns); -- see 3.6
 end record;
```

```

35 type Square(Side : Integer) is new
 Matrix_Rec(Rows => Side, Columns => Side);
36 type Double_Square(Number : Integer) is
 record
 Left : Square(Number);
 Right : Square(Number);
 end record;
37/2 task type Worker(Prio : System.Priority; Buf : access Buffer) is
 -- discriminants used to parameterize the task type (see 9.1)
 pragma Priority(Prio); -- see D.1
 entry Fill;
 entry Drain;
 end Worker;

```

### 3.7.1 Discriminant Constraints

- 1 A `discriminant_constraint` specifies the values of the discriminants for a given discriminated type.

*Syntax*

- 2 `discriminant_constraint ::=`  
`(discriminant_association {, discriminant_association})`
- 3 `discriminant_association ::=`  
`[discriminant_selector_name { | discriminant_selector_name} =>] expression`
- 4 A `discriminant_association` is said to be *named* if it has one or more `discriminant_selector_names`; it is otherwise said to be *positional*. In a `discriminant_constraint`, any positional associations shall precede any named associations.

*Name Resolution Rules*

- 5 Each `selector_name` of a named `discriminant_association` shall resolve to denote a discriminant of the subtype being constrained; the discriminants so named are the *associated discriminants* of the named association. For a positional association, the *associated discriminant* is the one whose `discriminant_specification` occurred in the corresponding position in the `known_discriminant_part` that defined the discriminants of the subtype being constrained.
- 6 The expected type for the `expression` in a `discriminant_association` is that of the associated discriminant(s).

*Legality Rules*

- 7/2 A `discriminant_constraint` is only allowed in a `subtype_indication` whose `subtype_mark` denotes either an unconstrained discriminated subtype, or an unconstrained access subtype whose designated subtype is an unconstrained discriminated subtype. However, in the case of an access subtype, a `discriminant_constraint` is illegal if the designated type has a partial view that is constrained or, for a general access subtype, has `default_expressions` for its discriminants. In addition to the places where Legality Rules normally apply (see 12.3), these rules apply also in the private part of an instance of a generic unit. In a generic body, this rule is checked presuming all formal access types of the generic might be general access types, and all untagged discriminated formal types of the generic might have `default_expressions` for their discriminants.
- 8 A named `discriminant_association` with more than one `selector_name` is allowed only if the named discriminants are all of the same type. A `discriminant_constraint` shall provide exactly one value for each discriminant of the subtype being constrained.

The expression associated with an access discriminant shall be of a type convertible to the anonymous access type. 9

#### Dynamic Semantics

A discriminant\_constraint is *compatible* with an unconstrained discriminated subtype if each discriminant value belongs to the subtype of the corresponding discriminant. 10

A composite value *satisfies* a discriminant constraint if and only if each discriminant of the composite value has the value imposed by the discriminant constraint. 11

For the elaboration of a discriminant\_constraint, the expressions in the discriminant\_associations are evaluated in an arbitrary order and converted to the type of the associated discriminant (which might raise Constraint\_Error — see 4.6); the expression of a named association is evaluated (and converted) once for each associated discriminant. The result of each evaluation and conversion is the value imposed by the constraint for the associated discriminant. 12

#### NOTES

56 The rules of the language ensure that a discriminant of an object always has a value, either from explicit or implicit initialization. 13

#### Examples

*Examples (using types declared above in clause 3.7):* 14

|                      |                                                                                       |
|----------------------|---------------------------------------------------------------------------------------|
| Large : Buffer(200); | <i>-- constrained, always 200 characters<br/>-- (explicit discriminant value)</i>     |
| Message : Buffer;    | <i>-- unconstrained, initially 100 characters<br/>-- (default discriminant value)</i> |
| Basis : Square(5);   | <i>-- constrained, always 5 by 5</i>                                                  |
| Illegal : Square;    | <i>-- illegal, a Square has to be constrained</i>                                     |

## 3.7.2 Operations of Discriminated Types

If a discriminated type has default\_expressions for its discriminants, then unconstrained variables of the type are permitted, and the discriminants of such a variable can be changed by assignment to the variable. For a formal parameter of such a type, an attribute is provided to determine whether the corresponding actual parameter is constrained or unconstrained. 1

#### Static Semantics

For a prefix A that is of a discriminated type (after any implicit dereference), the following attribute is defined: 2

A'Constrained

Yields the value True if A denotes a constant, a value, or a constrained variable, and False otherwise. 3

#### Erroneous Execution

The execution of a construct is erroneous if the construct has a constituent that is a name denoting a subcomponent that depends on discriminants, and the value of any of these discriminants is changed by this execution between evaluating the name and the last use (within this execution) of the subcomponent denoted by the name. 4

## 3.8 Record Types

- 1 A record object is a composite object consisting of named components. The value of a record object is a composite value consisting of the values of the components.

*Syntax*

```

2 record_type_definition ::= [[abstract] tagged] [limited] record_definition
3 record_definition ::=
4 record
5 component_list
6 end record
7 | null record
8 component_list ::=
9 component_item {component_item}
10 | {component_item} variant_part
11 | null;
12 component_item ::= component_declaration | aspect_clause
13 component_declaration ::=
14 defining_identifier_list : component_definition [: default_expression];

```

*Name Resolution Rules*

- 7 The expected type for the `default_expression`, if any, in a `component_declaration` is the type of the component.

*Legality Rules*

- 8/2 *This paragraph was deleted.*
- 9/2 Each `component_declaration` declares a component of the record type. Besides components declared by `component_declarations`, the components of a record type include any components declared by `discriminant_specifications` of the record type declaration. The identifiers of all components of a record type shall be distinct.
- 10 Within a `type_declaration`, a name that denotes a component, protected subprogram, or entry of the type is allowed only in the following cases:
- 11 • A name that denotes any component, protected subprogram, or entry is allowed within a representation item that occurs within the declaration of the composite type.
  - 12 • A name that denotes a noninherited discriminant is allowed within the declaration of the type, but not within the `discriminant_part`. If the discriminant is used to define the constraint of a component, the bounds of an entry family, or the constraint of the parent subtype in a `derived_type_definition` then its name shall appear alone as a `direct_name` (not as part of a larger expression or expanded name). A discriminant shall not be used to define the constraint of a scalar component.
- 13 If the name of the current instance of a type (see 8.6) is used to define the constraint of a component, then it shall appear as a `direct_name` that is the prefix of an `attribute_reference` whose result is of an access type, and the `attribute_reference` shall appear alone.

*Static Semantics*

If a `record_type_definition` includes the reserved word **limited**, the type is called an *explicitly limited record type*. 13.1/2

The `component_definition` of a `component_declaration` defines the (nominal) subtype of the component. 14  
If the reserved word **aliased** appears in the `component_definition`, then the component is aliased (see 3.10).

If the `component_list` of a record type is defined by the reserved word **null** and there are no discriminants, 15  
then the record type has no components and all records of the type are *null records*. A `record_definition` of **null record** is equivalent to `record null; end record`.

*Dynamic Semantics*

The elaboration of a `record_type_definition` creates the record type and its first subtype, and consists of 16  
the elaboration of the `record_definition`. The elaboration of a `record_definition` consists of the elaboration of its `component_list`, if any.

The elaboration of a `component_list` consists of the elaboration of the `component_items` and `variant_part`, 17  
if any, in the order in which they appear. The elaboration of a `component_declaration` consists of the elaboration of the `component_definition`.

Within the definition of a composite type, if a `component_definition` or `discrete_subtype_definition` (see 18/2  
9.5.2) includes a `name` that denotes a discriminant of the type, or that is an `attribute_reference` whose prefix denotes the current instance of the type, the expression containing the `name` is called a *per-object expression*, and the `constraint` or `range` being defined is called a *per-object constraint*. For the elaboration of a `component_definition` of a `component_declaration` or the `discrete_subtype_definition` of an `entry_declaration` for an entry family (see 9.5.2), if the component subtype is defined by an `access_definition` or if the `constraint` or `range` of the `subtype_indication` or `discrete_subtype_definition` is not a per-object constraint, then the `access_definition`, `subtype_indication`, or `discrete_subtype_definition` is elaborated. On the other hand, if the `constraint` or `range` is a per-object constraint, then the elaboration consists of the evaluation of any included expression that is not part of a per-object expression. Each such expression is evaluated once unless it is part of a named association in a discriminant constraint, in which case it is evaluated once for each associated discriminant.

When a per-object constraint is elaborated (as part of creating an object), each per-object expression of the constraint is evaluated. For other expressions, the values determined during the elaboration of the `component_definition` or `entry_declaration` are used. Any checks associated with the enclosing `subtype_indication` or `discrete_subtype_definition` are performed, including the subtype compatibility check (see 3.2.2), and the associated subtype is created. 18.1/1

## NOTES

57 A `component_declaration` with several identifiers is equivalent to a sequence of single `component_declarations`, as explained in 3.3.1. 19

58 The `default_expression` of a record component is only evaluated upon the creation of a default-initialized object of the record type (presuming the object has the component, if it is in a `variant_part` — see 3.3.1). 20

59 The subtype defined by a `component_definition` (see 3.6) has to be a definite subtype. 21

60 If a record type does not have a `variant_part`, then the same components are present in all values of the type. 22

61 A record type is limited if it has the reserved word **limited** in its definition, or if any of its components are limited (see 7.5). 23

62 The predefined operations of a record type include membership tests, qualification, and explicit conversion. If the record type is nonlimited, they also include assignment and the predefined equality operators. 24

- 25/2        63 A component of a record can be named with a `selected_component`. A value of a record can be specified with a `record_aggregate`.

*Examples*

26        *Examples of record type declarations:*

```
27 type Date is
 record
 Day : Integer range 1 .. 31;
 Month : Month_Name;
 Year : Integer range 0 .. 4000;
 end record;
28 type Complex is
 record
 Re : Real := 0.0;
 Im : Real := 0.0;
 end record;
```

29        *Examples of record variables:*

```
30 Tomorrow, Yesterday : Date;
 A, B, C : Complex;
31 -- both components of A, B, and C are implicitly initialized to zero
```

### 3.8.1 Variant Parts and Discrete Choices

- 1        A record type with a `variant_part` specifies alternative lists of components. Each `variant` defines the components for the value or values of the discriminant covered by its `discrete_choice_list`.

*Syntax*

```
2 variant_part ::=
 case discriminant_direct_name is
 variant
 {variant}
 end case;
3 variant ::=
 when discrete_choice_list =>
 component_list
4 discrete_choice_list ::= discrete_choice { | discrete_choice}
5 discrete_choice ::= expression | discrete_range | others
```

*Name Resolution Rules*

- 6        The `discriminant_direct_name` shall resolve to denote a discriminant (called the *discriminant of the variant\_part*) specified in the `known_discriminant_part` of the `full_type_declaration` that contains the `variant_part`. The expected type for each `discrete_choice` in a `variant` is the type of the discriminant of the `variant_part`.

*Legality Rules*

- 7        The discriminant of the `variant_part` shall be of a discrete type.
- 8        The expressions and `discrete_ranges` given as `discrete_choices` in a `variant_part` shall be static. The `discrete_choice others` shall appear alone in a `discrete_choice_list`, and such a `discrete_choice_list`, if it appears, shall be the last one in the enclosing construct.

A discrete\_choice is defined to *cover a value* in the following cases:

- A discrete\_choice that is an expression covers a value if the value equals the value of the expression converted to the expected type.
- A discrete\_choice that is a discrete\_range covers all values (possibly none) that belong to the range.
- The discrete\_choice others covers all values of its expected type that are not covered by previous discrete\_choice\_lists of the same construct.

A discrete\_choice\_list covers a value if one of its discrete\_choices covers the value.

The possible values of the discriminant of a variant\_part shall be covered as follows:

- If the discriminant is of a static constrained scalar subtype, then each non-others discrete\_choice shall cover only values in that subtype, and each value of that subtype shall be covered by some discrete\_choice (either explicitly or by others);
- If the type of the discriminant is a descendant of a generic formal scalar type then the variant\_part shall have an others discrete\_choice;
- Otherwise, each value of the base range of the type of the discriminant shall be covered (either explicitly or by others).

Two distinct discrete\_choices of a variant\_part shall not cover the same value.

#### *Static Semantics*

If the component\_list of a variant is specified by null, the variant has no components.

The discriminant of a variant\_part is said to *govern* the variant\_part and its variants. In addition, the discriminant of a derived type governs a variant\_part and its variants if it corresponds (see 3.7) to the discriminant of the variant\_part.

#### *Dynamic Semantics*

A record value contains the values of the components of a particular variant only if the value of the discriminant governing the variant is covered by the discrete\_choice\_list of the variant. This rule applies in turn to any further variant that is, itself, included in the component\_list of the given variant.

The elaboration of a variant\_part consists of the elaboration of the component\_list of each variant in the order in which they appear.

#### *Examples*

*Example of record type with a variant part:*

```
type Device is (Printer, Disk, Drum);
type State is (Open, Closed);
type Peripheral(Unit : Device := Disk) is
 record
 Status : State;
 case Unit is
 when Printer =>
 Line_Count : Integer range 1 .. Page_Size;
 when others =>
 Cylinder : Cylinder_Index;
 Track : Track_Number;
 end case;
 end record;
```

26     *Examples of record subtypes:*

27       **subtype** Drum\_Unit **is** Peripheral(Drum);  
          **subtype** Disk\_Unit **is** Peripheral(Disk);

28     *Examples of constrained record variables:*

29       Writer    : Peripheral(Unit  => Printer);  
          Archive  : Disk\_Unit;

## 3.9 Tagged Types and Type Extensions

- 1     Tagged types and type extensions support object-oriented programming, based on inheritance with extension and run-time polymorphism via *dispatching operations*.

### Static Semantics

- 2/2   A record type or private type that has the reserved word **tagged** in its declaration is called a *tagged type*. In addition, an interface type is a tagged type, as is a task or protected type derived from an interface (see 3.9.4). When deriving from a tagged type, as for any derived type, additional primitive subprograms may be defined, and inherited primitive subprograms may be overridden. The derived type is called an *extension* of its ancestor types, or simply a *type extension*.
- 2.1/2 Every type extension is also a tagged type, and is a *record extension* or a *private extension* of some other tagged type, or a non-interface synchronized tagged type (see 3.9.4). A record extension is defined by a **derived\_type\_definition** with a **record\_extension\_part** (see 3.9.1), which may include the definition of additional components. A private extension, which is a partial view of a record extension or of a synchronized tagged type, can be declared in the visible part of a package (see 7.3) or in a generic formal part (see 12.5.1).
- 3     An object of a tagged type has an associated (run-time) *tag* that identifies the specific tagged type used to create the object originally. The tag of an operand of a class-wide tagged type *TClass* controls which subprogram body is to be executed when a primitive subprogram of type *T* is applied to the operand (see 3.9.2); using a tag to control which body to execute is called *dispatching*.
- 4/2   The tag of a specific tagged type identifies the **full\_type\_declaration** of the type, and for a type extension, is sufficient to uniquely identify the type among all descendants of the same ancestor. If a declaration for a tagged type occurs within a **generic\_package\_declaration**, then the corresponding type declarations in distinct instances of the generic package are associated with distinct tags. For a tagged type that is local to a generic package body and with all of its ancestors (if any) also local to the generic body, the language does not specify whether repeated instantiations of the generic body result in distinct tags.
- 5     The following language-defined library package exists:

```
6/2 package Ada.Tags is
 pragma Preelaborate(Tags);
 type Tag is private;
 pragma Preelaborable_Initialization(Tag);
 No_Tag : constant Tag;
 function Expanded_Name(T : Tag) return String;
 function Wide_Expanded_Name(T : Tag) return Wide_String;
 function Wide_Wide_Expanded_Name(T : Tag) return Wide_Wide_String;
 function External_Tag(T : Tag) return String;
 function Internal_Tag(External : String) return Tag;
 function Descendant_Tag(External : String; Ancestor : Tag) return Tag;
 function Is_Descendant_At_Same_Level(Descendant, Ancestor : Tag)
 return Boolean;
```

```

function Parent_Tag (T : Tag) return Tag; 7.2/2
type Tag_Array is array (Positive range <>) of Tag; 7.3/2
function Interface_Ancestor_Tags (T : Tag) return Tag_Array; 7.4/2
 Tag_Error : exception; 8
private 9
 . . . -- not specified by the language
end Ada.Tags;

```

No\_Tag is the default initial value of type Tag. 9.1/2

The function Wide\_Wide\_Expanded\_Name returns the full expanded name of the first subtype of the specific type identified by the tag, in upper case, starting with a root library unit. The result is implementation defined if the type is declared within an unnamed block\_statement. 10/2

The function Expanded\_Name (respectively, Wide\_Expanded\_Name) returns the same sequence of graphic characters as that defined for Wide\_Wide\_Expanded\_Name, if all the graphic characters are defined in Character (respectively, Wide\_Character); otherwise, the sequence of characters is implementation defined, but no shorter than that returned by Wide\_Wide\_Expanded\_Name for the same value of the argument. 10.1/2

The function External\_Tag returns a string to be used in an external representation for the given tag. The call External\_Tag(S'Tag) is equivalent to the attribute\_reference S'External\_Tag (see 13.3). 11

The string returned by the functions Expanded\_Name, Wide\_Expanded\_Name, Wide\_Wide\_Expanded\_Name, and External\_Tag has lower bound 1. 11.1/2

The function Internal\_Tag returns a tag that corresponds to the given external tag, or raises Tag\_Error if the given string is not the external tag for any specific type of the partition. Tag\_Error is also raised if the specific type identified is a library-level type whose tag has not yet been created (see 13.14). 12/2

The function Descendant\_Tag returns the (internal) tag for the type that corresponds to the given external tag and is both a descendant of the type identified by the Ancestor tag and has the same accessibility level as the identified ancestor. Tag\_Error is raised if External is not the external tag for such a type. Tag\_Error is also raised if the specific type identified is a library-level type whose tag has not yet been created. 12.1/2

The function Is\_Descendant\_At\_Same\_Level returns True if the Descendant tag identifies a type that is both a descendant of the type identified by Ancestor and at the same accessibility level. If not, it returns False. 12.2/2

The function Parent\_Tag returns the tag of the parent type of the type whose tag is T. If the type does not have a parent type (that is, it was not declared by a derived\_type\_declaration), then No\_Tag is returned. 12.3/2

The function Interface\_Ancestor\_Tags returns an array containing the tag of each interface ancestor type of the type whose tag is T, other than T itself. The lower bound of the returned array is 1, and the order of the returned tags is unspecified. Each tag appears in the result exactly once. If the type whose tag is T has no interface ancestors, a null array is returned. 12.4/2

For every subtype S of a tagged type T (specific or class-wide), the following attributes are defined: 13

S'Class 14  
S'Class denotes a subtype of the class-wide type (called TClass in this International Standard) for the class rooted at T (or if S already denotes a class-wide subtype, then S'Class is the same as S).

S'Class is unconstrained. However, if S is constrained, then the values of S'Class are only those that when converted to the type T belong to S. 15

- 16 S'Tag        S'Tag denotes the tag of the type *T* (or if *T* is class-wide, the tag of the root type of the corresponding class). The value of this attribute is of type Tag.
- 17 Given a prefix X that is of a class-wide tagged type (after any implicit dereference), the following attribute is defined:
- 18 X'Tag        X'Tag denotes the tag of X. The value of this attribute is of type Tag.

18.1/2 The following language-defined generic function exists:

```
18.2/2 generic
 type T (<>) is abstract tagged limited private;
 type Parameters (<>) is limited private;
 with function Constructor (Params : not null access Parameters)
 return T is abstract;
 function Ada.Tags.Generic_Dispatching_Constructor
 (The_Tag : Tag;
 Params : not null access Parameters) return T'Class;
 pragma Preelaborate(Generic_Dispatching_Constructor);
 pragma Convention(Intrinsic, Generic_Dispatching_Constructor);
```

18.3/2 Tags.Generic\_Dispatching\_Constructor provides a mechanism to create an object of an appropriate type from just a tag value. The function Constructor is expected to create the object given a reference to an object of type Parameters.

#### *Dynamic Semantics*

- 19 The tag associated with an object of a tagged type is determined as follows:
- 20 • The tag of a stand-alone object, a component, or an aggregate of a specific tagged type *T* identifies *T*.
  - 21 • The tag of an object created by an allocator for an access type with a specific designated tagged type *T*, identifies *T*.
  - 22 • The tag of an object of a class-wide tagged type is that of its initialization expression.
  - 23 • The tag of the result returned by a function whose result type is a specific tagged type *T* identifies *T*.
  - 24/2 • The tag of the result returned by a function with a class-wide result type is that of the return object.
- 25 The tag is preserved by type conversion and by parameter passing. The tag of a value is the tag of the associated object (see 6.2).
- 25.1/2 Tag\_Error is raised by a call of Descendant\_Tag, Expanded\_Name, External\_Tag, Interface\_Ancestor\_Tag, Is\_Descendant\_At\_Same\_Level, or Parent\_Tag if any tag passed is No\_Tag.
- 25.2/2 An instance of Tags.Generic\_Dispatching\_Constructor raises Tag\_Error if The\_Tag does not represent a concrete descendant of *T* or if the innermost master (see 7.6.1) of this descendant is not also a master of the instance. Otherwise, it dispatches to the primitive function denoted by the formal Constructor for the type identified by The\_Tag, passing Params, and returns the result. Any exception raised by the function is propagated.

#### *Erroneous Execution*

- 25.3/2 If an internal tag provided to an instance of Tags.Generic\_Dispatching\_Constructor or to any subprogram declared in package Tags identifies either a type that is not library-level and whose tag has not been created (see 13.14), or a type that does not exist in the partition at the time of the call, then execution is erroneous.

*Implementation Permissions*

The implementation of Internal\_Tag and Descendant\_Tag may raise Tag\_Error if no specific type corresponding to the string External passed as a parameter exists in the partition at the time the function is called, or if there is no such type whose innermost master is a master of the point of the function call. 26/2

*Implementation Advice*

Internal\_Tag should return the tag of a type whose innermost master is the master of the point of the function call. 26.1/2

## NOTES

64 A type declared with the reserved word **tagged** should normally be declared in a package\_specification, so that new primitive subprograms can be declared for it. 27

65 Once an object has been created, its tag never changes. 28

66 Class-wide types are defined to have unknown discriminants (see 3.7). This means that objects of a class-wide type have to be explicitly initialized (whether created by an object\_declaration or an allocator), and that aggregates have to be explicitly qualified with a specific type when their expected type is class-wide. 29

*This paragraph was deleted.* 30/2

67 The capability provided by Tags.Generic\_Dispatching\_Constructor is sometimes known as a *factory*. 30.1/2

*Examples**Examples of tagged record types:*

```
type Point is tagged
 record
 X, Y : Real := 0.0;
 end record;

type Expression is tagged null record;
 -- Components will be added by each extension
```

**3.9.1 Type Extensions**

Every type extension is a tagged type, and is a *record extension* or a *private extension* of some other tagged type, or a non-interface synchronized tagged type.. 1/2

*Syntax*

record\_extension\_part ::= with record\_definition 2

*Legality Rules*

The parent type of a record extension shall not be a class-wide type nor shall it be a synchronized tagged type (see 3.9.4). If the parent type or any progenitor is nonlimited, then each of the components of the record\_extension\_part shall be nonlimited. In addition to the places where Legality Rules normally apply (see 12.3), these rules apply also in the private part of an instance of a generic unit. 3/2

Within the body of a generic unit, or the body of any of its descendant library units, a tagged type shall not be declared as a descendant of a formal type declared within the formal part of the generic unit. 4/2

*Static Semantics*

A record extension is a *null extension* if its declaration has no known\_discriminant\_part and its record\_extension\_part includes no component\_declarations. 4.1/2

*Dynamic Semantics*

The elaboration of a record\_extension\_part consists of the elaboration of the record\_definition. 5

## NOTES

- 6        68 The term “type extension” refers to a type as a whole. The term “extension part” refers to the piece of text that defines the additional components (if any) the type extension has relative to its specified ancestor type.
- 7/2      69 When an extension is declared immediately within a body, primitive subprograms are inherited and are overridable, but new primitive subprograms cannot be added.
- 8        70 A name that denotes a component (including a discriminant) of the parent type is not allowed within the `record_extension_part`. Similarly, a name that denotes a component defined within the `record_extension_part` is not allowed within the `record_extension_part`. It is permissible to use a name that denotes a discriminant of the record extension, providing there is a new `known_discriminant_part` in the enclosing type declaration. (The full rule is given in 3.8.)
- 9        71 Each visible component of a record extension has to have a unique name, whether the component is (visibly) inherited from the parent type or declared in the `record_extension_part` (see 8.3).

*Examples*

10      Examples of record extensions (of types defined above in 3.9):

```

11 type Painted_Point is new Point with
 record
 Paint : Color := White;
 end record;
 -- Components X and Y are inherited
12 Origin : constant Painted_Point := (X | Y => 0.0, Paint => Black);
13 type Literal is new Expression with
 record
 Value : Real;
 end record;
14 type Expr_Ptr is access all Expression'Class;
 -- see 3.10
15 type Binary_Operation is new Expression with
 record
 -- an internal node in an Expression tree
 Left, Right : Expr_Ptr;
 end record;
16 type Addition is new Binary_Operation with null record;
 type Subtraction is new Binary_Operation with null record;
 -- No additional components needed for these extensions
17 Tree : Expr_Ptr := -- A tree representation of "5.0 + (13.0-7.0)"
 new Addition'(
 Left => new Literal'(Value => 5.0),
 Right => new Subtraction'(
 Left => new Literal'(Value => 13.0),
 Right => new Literal'(Value => 7.0)));

```

### 3.9.2 Dispatching Operations of Tagged Types

- 1/2      The primitive subprograms of a tagged type, the subprograms declared by `formal_abstract_subprogram_declarations`, and the stream attributes of a specific tagged type that are available (see 13.13.2) at the end of the declaration list where the type is declared are called *dispatching operations*. A dispatching operation can be called using a statically determined *controlling tag*, in which case the body to be executed is determined at compile time. Alternatively, the controlling tag can be dynamically determined, in which case the call *dispatches* to a body that is determined at run time; such a call is termed a *dispatching call*. As explained below, the properties of the operands and the context of a particular call on a dispatching operation determine how the controlling tag is determined, and hence whether or not the call is a dispatching call. Run-time polymorphism is achieved when a dispatching operation is called by a dispatching call.

*Static Semantics*

A *call on a dispatching operation* is a call whose name or prefix denotes the declaration of a dispatching operation. A *controlling operand* in a call on a dispatching operation of a tagged type  $T$  is one whose corresponding formal parameter is of type  $T$  or is of an anonymous access type with designated type  $T$ ; the corresponding formal parameter is called a *controlling formal parameter*. If the controlling formal parameter is an access parameter, the controlling operand is the object designated by the actual parameter, rather than the actual parameter itself. If the call is to a (primitive) function with result type  $T$ , then the call has a *controlling result* — the context of the call can control the dispatching. Similarly, if the call is to a function with access result type designating  $T$ , then the call has a *controlling access result*, and the context can similarly control dispatching.

A name or expression of a tagged type is either *statically tagged*, *dynamically tagged*, or *tag indeterminate*, according to whether, when used as a controlling operand, the tag that controls dispatching is determined statically by the operand's (specific) type, dynamically by its tag at run time, or from context. A *qualified\_expression* or parenthesized expression is statically, dynamically, or indeterminately tagged according to its operand. For other kinds of names and expressions, this is determined as follows:

- The name or expression is *statically tagged* if it is of a specific tagged type and, if it is a call with a controlling result or controlling access result, it has at least one statically tagged controlling operand;
- The name or expression is *dynamically tagged* if it is of a class-wide type, or it is a call with a controlling result or controlling access result and at least one dynamically tagged controlling operand;
- The name or expression is *tag indeterminate* if it is a call with a controlling result or controlling access result, all of whose controlling operands (if any) are tag indeterminate.

A *type\_conversion* is statically or dynamically tagged according to whether the type determined by the *subtype\_mark* is specific or class-wide, respectively. For an object that is designated by an expression whose expected type is an anonymous access-to-specific tagged type, the object is dynamically tagged if the expression, ignoring enclosing parentheses, is of the form  $X'Access$ , where  $X$  is of a class-wide type, or is of the form  $\text{new } T'(..)$ , where  $T$  denotes a class-wide subtype. Otherwise, the object is statically or dynamically tagged according to whether the designated type of the type of the expression is specific or class-wide, respectively.

*Legality Rules*

A call on a dispatching operation shall not have both dynamically tagged and statically tagged controlling operands.

If the expected type for an expression or name is some specific tagged type, then the expression or name shall not be dynamically tagged unless it is a controlling operand in a call on a dispatching operation. Similarly, if the expected type for an expression is an anonymous access-to-specific tagged type, then the object designated by the expression shall not be dynamically tagged unless it is a controlling operand in a call on a dispatching operation.

In the declaration of a dispatching operation of a tagged type, everywhere a subtype of the tagged type appears as a subtype of the profile (see 6.1), it shall statically match the first subtype of the tagged type. If the dispatching operation overrides an inherited subprogram, it shall be subtype conformant with the inherited subprogram. The convention of an inherited dispatching operation is the convention of the corresponding primitive operation of the parent or progenitor type. The default convention of a dispatching operation that overrides an inherited primitive operation is the convention of the inherited operation; if the

operation overrides multiple inherited operations, then they shall all have the same convention. An explicitly declared dispatching operation shall not be of convention Intrinsic.

- 11/2 The `default_expression` for a controlling formal parameter of a dispatching operation shall be tag indeterminate.
- 11.1/2 If a dispatching operation is defined by a `subprogram_renaming_declaration` or the instantiation of a generic subprogram, any access parameter of the renamed subprogram or the generic subprogram that corresponds to a controlling access parameter of the dispatching operation, shall have a subtype that excludes null.
- 12 A given subprogram shall not be a dispatching operation of two or more distinct tagged types.
- 13 The explicit declaration of a primitive subprogram of a tagged type shall occur before the type is frozen (see 13.14). For example, new dispatching operations cannot be added after objects or values of the type exist, nor after deriving a record extension from it, nor after a body.

#### *Dynamic Semantics*

- 14 For the execution of a call on a dispatching operation of a type  $T$ , the *controlling tag value* determines which subprogram body is executed. The controlling tag value is defined as follows:
  - 15 • If one or more controlling operands are statically tagged, then the controlling tag value is *statically determined* to be the tag of  $T$ .
  - 16 • If one or more controlling operands are dynamically tagged, then the controlling tag value is not statically determined, but is rather determined by the tags of the controlling operands. If there is more than one dynamically tagged controlling operand, a check is made that they all have the same tag. If this check fails, `Constraint_Error` is raised unless the call is a `function_call` whose `name` denotes the declaration of an equality operator (predefined or user defined) that returns Boolean, in which case the result of the call is defined to indicate inequality, and no `subprogram_body` is executed. This check is performed prior to evaluating any tag-indeterminate controlling operands.
  - 17/2 • If all of the controlling operands (if any) are tag-indeterminate, then:
    - 18/2 • If the call has a controlling result or controlling access result and is itself, or designates, a (possibly parenthesized or qualified) controlling operand of an enclosing call on a dispatching operation of a descendant of type  $T$ , then its controlling tag value is determined by the controlling tag value of this enclosing call;
    - 18.1/2 • If the call has a controlling result or controlling access result and (possibly parenthesized, qualified, or dereferenced) is the expression of an `assignment_statement` whose target is of a class-wide type, then its controlling tag value is determined by the target;
    - 19 • Otherwise, the controlling tag value is statically determined to be the tag of type  $T$ .
  - 20/2 For the execution of a call on a dispatching operation, the action performed is determined by the properties of the corresponding dispatching operation of the specific type identified by the controlling tag value. If the corresponding operation is explicitly declared for this type, even if the declaration occurs in a private part, then the action comprises an invocation of the explicit body for the operation. If the corresponding operation is implicitly declared for this type:
    - 20.1/2 • if the operation is implemented by an entry or protected subprogram (see 9.1 and 9.4), then the action comprises a call on this entry or protected subprogram, with the target object being given by the first actual parameter of the call, and the actual parameters of the entry or protected subprogram being given by the remaining actual parameters of the call, if any;
    - 20.2/2 • otherwise, the action is the same as the action for the corresponding operation of the parent type.

## NOTES

72 The body to be executed for a call on a dispatching operation is determined by the tag; it does not matter whether that tag is determined statically or dynamically, and it does not matter whether the subprogram's declaration is visible at the place of the call. 21

73 This subclause covers calls on dispatching subprograms of a tagged type. Rules for tagged type membership tests are described in 4.5.2. Controlling tag determination for an `assignment_statement` is described in 5.2. 22/2

74 A dispatching call can dispatch to a body whose declaration is not visible at the place of the call. 23

75 A call through an access-to-subprogram value is never a dispatching call, even if the access value designates a dispatching operation. Similarly a call whose prefix denotes a `subprogram_renaming_declaration` cannot be a dispatching call unless the renaming itself is the declaration of a primitive subprogram. 24

### 3.9.3 Abstract Types and Subprograms

An *abstract type* is a tagged type intended for use as an ancestor of other types, but which is not allowed to have objects of its own. An *abstract subprogram* is a subprogram that has no body, but is intended to be overridden at some point when inherited. Because objects of an abstract type cannot be created, a dispatching call to an abstract subprogram always dispatches to some overriding body. 1/2

#### *Syntax*

`abstract_subprogram_declaration ::=`  
`[overriding_indicator]`  
`subprogram_specification is abstract;`

1.1/2

#### *Static Semantics*

Interface types (see 3.9.4) are abstract types. In addition, a tagged type that has the reserved word **abstract** in its declaration is an abstract type. The class-wide type (see 3.4.1) rooted at an abstract type is not itself an abstract type. 1.2/2

#### *Legality Rules*

Only a tagged type shall have the reserved word **abstract** in its declaration. 2/2

A subprogram declared by an `abstract_subprogram_declaration` or a `formal_abstract_subprogram_declaration` (see 12.6) is an *abstract subprogram*. If it is a primitive subprogram of a tagged type, then the tagged type shall be abstract. 3/2

If a type has an implicitly declared primitive subprogram that is inherited or is the predefined equality operator, and the corresponding primitive subprogram of the parent or ancestor type is abstract or is a function with a controlling access result, or if a type other than a null extension inherits a function with a controlling result, then: 4/2

- If the type is abstract or untagged, the implicitly declared subprogram is *abstract*. 5/2
- Otherwise, the subprogram shall be overridden with a nonabstract subprogram or, in the case of a private extension inheriting a function with a controlling result, have a full type that is a null extension; for a type declared in the visible part of a package, the overriding may be either in the visible or the private part. Such a subprogram is said to *require overriding*. However, if the type is a generic formal type, the subprogram need not be overridden for the formal type itself; a nonabstract version will necessarily be provided by the actual type. 6/2

A call on an abstract subprogram shall be a dispatching call; nondispatching calls to an abstract subprogram are not allowed. 7

- 8 The type of an **aggregate**, or of an object created by an **object\_declaration** or an **allocator**, or a generic formal object of mode **in**, shall not be abstract. The type of the target of an assignment operation (see 5.2) shall not be abstract. The type of a component shall not be abstract. If the result type of a function is abstract, then the function shall be abstract.
- 9 If a partial view is not abstract, the corresponding full view shall not be abstract. If a generic formal type is abstract, then for each primitive subprogram of the formal that is not abstract, the corresponding primitive subprogram of the actual shall not be abstract.
- 10 For an abstract type declared in a visible part, an abstract primitive subprogram shall not be declared in the private part, unless it is overriding an abstract subprogram implicitly declared in the visible part. For a tagged type declared in a visible part, a primitive function with a controlling result shall not be declared in the private part, unless it is overriding a function implicitly declared in the visible part.
- 11/2 A generic actual subprogram shall not be an abstract subprogram unless the generic formal subprogram is declared by a **formal\_abstract\_subprogram\_declaration**. The prefix of an **attribute\_reference** for the **Access**, **Unchecked\_Access**, or **Address** attributes shall not denote an abstract subprogram.

*Dynamic Semantics*

- 11.1/2 The elaboration of an **abstract\_subprogram\_declaration** has no effect.

## NOTES

- 12 76 Abstractness is not inherited; to declare an abstract type, the reserved word **abstract** has to be used in the declaration of the type extension.
- 13 77 A class-wide type is never abstract. Even if a class is rooted at an abstract type, the class-wide type for the class is not abstract, and an object of the class-wide type can be created; the tag of such an object will identify some nonabstract type in the class.

*Examples*

- 14 Example of an abstract type representing a set of natural numbers:

```
15 package Sets is
 subtype Element_Type is Natural;
 type Set is abstract tagged null record;
 function Empty return Set is abstract;
 function Union(Left, Right : Set) return Set is abstract;
 function Intersection(Left, Right : Set) return Set is abstract;
 function Unit_Set(Element : Element_Type) return Set is abstract;
 procedure Take(Element : out Element_Type;
 From : in out Set) is abstract;
end Sets;
```

## NOTES

- 16 78 Notes on the example: Given the above abstract type, one could then derive various (nonabstract) extensions of the type, representing alternative implementations of a set. One might use a bit vector, but impose an upper bound on the largest element representable, while another might use a hash table, trading off space for flexibility.

### 3.9.4 Interface Types

- 1/2 An interface type is an abstract tagged type that provides a restricted form of multiple inheritance. A tagged type, task type, or protected type may have one or more interface types as ancestors.

*Syntax*

- 2/2 **interface\_type\_definition ::=**  
[**limited** | **task** | **protected** | **synchronized**] **interface** [**and** **interface\_list**]
- 3/2 **interface\_list ::=** *interface\_subtype\_mark* {**and** *interface\_subtype\_mark*}

*Static Semantics*

An interface type (also called an *interface*) is a specific abstract tagged type that is defined by an `interface_type_definition`. 4/2

An interface with the reserved word **limited**, **task**, **protected**, or **synchronized** in its definition is termed, respectively, a *limited interface*, a *task interface*, a *protected interface*, or a *synchronized interface*. In addition, all task and protected interfaces are synchronized interfaces, and all synchronized interfaces are limited interfaces. 5/2

A task or protected type derived from an interface is a tagged type. Such a tagged type is called a *synchronized tagged type*, as are synchronized interfaces and private extensions whose declaration includes the reserved word **synchronized**. 6/2

A task interface is an abstract task type. A protected interface is an abstract protected type. 7/2

An interface type has no components. 8/2

An `interface_subtype_mark` in an `interface_list` names a *progenitor subtype*; its type is the *progenitor type*. An interface type inherits user-defined primitive subprograms from each progenitor type in the same way that a derived type inherits user-defined primitive subprograms from its progenitor types (see 3.4). 9/2

*Legality Rules*

All user-defined primitive subprograms of an interface type shall be abstract subprograms or null procedures. 10/2

The type of a subtype named in an `interface_list` shall be an interface type. 11/2

A type derived from a nonlimited interface shall be nonlimited. 12/2

An interface derived from a task interface shall include the reserved word **task** in its definition; any other type derived from a task interface shall be a private extension or a task type declared by a task declaration (see 9.1). 13/2

An interface derived from a protected interface shall include the reserved word **protected** in its definition; any other type derived from a protected interface shall be a private extension or a protected type declared by a protected declaration (see 9.4). 14/2

An interface derived from a synchronized interface shall include one of the reserved words **task**, **protected**, or **synchronized** in its definition; any other type derived from a synchronized interface shall be a private extension, a task type declared by a task declaration, or a protected type declared by a protected declaration. 15/2

No type shall be derived from both a task interface and a protected interface. 16/2

In addition to the places where Legality Rules normally apply (see 12.3), these rules apply also in the private part of an instance of a generic unit. 17/2

*Dynamic Semantics*

The elaboration of an `interface_type_definition` has no effect. 18/2

## NOTES

79 Nonlimited interface types have predefined nonabstract equality operators. These may be overridden with user-defined abstract equality operators. Such operators will then require an explicit overriding for any nonabstract descendant of the interface. 19/2

*Examples*

20/2 Example of a limited interface and a synchronized interface extending it:

```

21/2 type Queue is limited interface;
procedure Append(Q : in out Queue; Person : in Person_Name) is abstract;
procedure Remove_First(Q : in out Queue;
 Person : out Person_Name) is abstract;
function Cur_Count(Q : in Queue) return Natural is abstract;
function Max_Count(Q : in Queue) return Natural is abstract;
-- See 3.10.1 for Person_Name.

22/2 Queue_Error : exception;
-- Append raises Queue_Error if Count(Q) = Max_Count(Q)
-- Remove_First raises Queue_Error if Count(Q) = 0

23/2 type Synchronized_Queue is synchronized interface and Queue; -- see 9.11
procedure Append_Wait(Q : in out Synchronized_Queue;
 Person : in Person_Name) is abstract;
procedure Remove_First_Wait(Q : in out Synchronized_Queue;
 Person : out Person_Name) is abstract;

24/2 ...
25/2 procedure Transfer(From : in out Queue'Class;
 To : in out Queue'Class;
 Number : in Natural := 1) is
 Person : Person_Name;
begin
 for I in 1..Number loop
 Remove_First(From, Person);
 Append(To, Person);
 end loop;
end Transfer;

```

26/2 This defines a Queue interface defining a queue of people. (A similar design could be created to define any kind of queue simply by replacing Person\_Name by an appropriate type.) The Queue interface has four dispatching operations, Append, Remove\_First, Cur\_Count, and Max\_Count. The body of a class-wide operation, Transfer is also shown. Every non-abstract extension of Queue must provide implementations for at least its four dispatching operations, as they are abstract. Any object of a type derived from Queue may be passed to Transfer as either the From or the To operand. The two operands need not be of the same type in any given call.

27/2 The Synchronized\_Queue interface inherits the four dispatching operations from Queue and adds two additional dispatching operations, which wait if necessary rather than raising the Queue\_Error exception. This synchronized interface may only be implemented by a task or protected type, and as such ensures safe concurrent access.

28/2 Example use of the interface:

```

29/2 type Fast_Food_Queue is new Queue with record ...;
procedure Append(Q : in out Fast_Food_Queue; Person : in Person_Name);
procedure Remove_First(Q : in out Fast_Food_Queue; Person : in Person_Name);
function Cur_Count(Q : in Fast_Food_Queue) return Natural;
function Max_Count(Q : in Fast_Food_Queue) return Natural;

30/2 ...
31/2 Cashier, Counter : Fast_Food_Queue;
32/2 ...
-- Add George (see 3.10.1) to the cashier's queue:
Append (Cashier, George);
-- After payment, move George to the sandwich counter queue:
Transfer (Cashier, Counter);
...
```

An interface such as Queue can be used directly as the parent of a new type (as shown here), or can be used as a progenitor when a type is derived. In either case, the primitive operations of the interface are inherited. For Queue, the implementation of the four inherited routines must be provided. Inside the call of Transfer, calls will dispatch to the implementations of Append and Remove\_First for type Fast\_Food\_Queue.

33/2

*Example of a task interface:*

34/2

```
type Serial_Device is task interface; -- see 9.1
procedure Read (Dev : in Serial_Device; C : out Character) is abstract;
procedure Write(Dev : in Serial_Device; C : in Character) is abstract;
```

35/2

The Serial\_Device interface has two dispatching operations which are intended to be implemented by task entries (see 9.1).

36/2

## 3.10 Access Types

A value of an access type (an *access value*) provides indirect access to the object or subprogram it designates. Depending on its type, an access value can designate either subprograms, objects created by allocators (see 4.8), or more generally aliased objects of an appropriate type.

1

### Syntax

```
access_type_definition ::= [null_exclusion] access_to_object_definition
| [null_exclusion] access_to_subprogram_definition
access_to_object_definition ::= access [general_access_modifier] subtype_indication
general_access_modifier ::= all | constant
access_to_subprogram_definition ::= access [protected] procedure parameter_profile
| access [protected] function parameter_and_result_profile
null_exclusion ::= not null
access_definition ::= [null_exclusion] access [constant] subtype_mark
| [null_exclusion] access [protected] procedure parameter_profile
| [null_exclusion] access [protected] function parameter_and_result_profile
```

2/2

3

4

5

5.1/2

6/2

### Static Semantics

There are two kinds of access types, *access-to-object* types, whose values designate objects, and *access-to-subprogram* types, whose values designate subprograms. Associated with an access-to-object type is a *storage pool*; several access types may share the same storage pool. All descendants of an access type share the same storage pool. A storage pool is an area of storage used to hold dynamically allocated objects (called *pool elements*) created by allocators; storage pools are described further in 13.11, “Storage Management”.

7/1

Access-to-object types are further subdivided into *pool-specific* access types, whose values can designate only the elements of their associated storage pool, and *general* access types, whose values can designate the elements of any storage pool, as well as aliased objects created by declarations rather than allocators, and aliased subcomponents of other objects.

8

- 9/2 A view of an object is defined to be *aliased* if it is defined by an `object_declaration` or `component_definition` with the reserved word **aliased**, or by a renaming of an aliased view. In addition, the dereference of an access-to-object value denotes an aliased view, as does a view conversion (see 4.6) of an aliased view. The current instance of a limited tagged type, a protected type, a task type, or a type that has the reserved word **limited** in its full definition is also defined to be aliased. Finally, a formal parameter or generic formal object of a tagged type is defined to be aliased. Aliased views are the ones that can be designated by an access value.
- 10 An `access_to_object_definition` defines an access-to-object type and its first subtype; the `subtype_indication` defines the *designated subtype* of the access type. If a `general_access_modifier` appears, then the access type is a general access type. If the modifier is the reserved word **constant**, then the type is an *access-to-constant type*; a designated object cannot be updated through a value of such a type. If the modifier is the reserved word **all**, then the type is an *access-to-variable type*; a designated object can be both read and updated through a value of such a type. If no `general_access_modifier` appears in the `access_to_object_definition`, the access type is a pool-specific access-to-variable type.
- 11 An `access_to_subprogram_definition` defines an access-to-subprogram type and its first subtype; the `parameter_profile` or `parameter_and_result_profile` defines the *designated profile* of the access type. There is a *calling convention* associated with the designated profile; only subprograms with this calling convention can be designated by values of the access type. By default, the calling convention is “*protected*” if the reserved word **protected** appears, and “Ada” otherwise. See Annex B for how to override this default.
- 12/2 An `access_definition` defines an anonymous general access type or an anonymous access-to-subprogram type. For a general access type, the `subtype_mark` denotes its *designated subtype*; if the `general_access_modifier constant` appears, the type is an access-to-constant type; otherwise it is an access-to-variable type. For an access-to-subprogram type, the `parameter_profile` or `parameter_and_result_profile` denotes its *designated profile*.
- 13/2 For each access type, there is a null access value designating no entity at all, which can be obtained by (implicitly) converting the literal **null** to the access type. The null value of an access type is the default initial value of the type. Non-null values of an access-to-object type are obtained by evaluating an `allocator`, which returns an access value designating a newly created object (see 3.10.2), or in the case of a general access-to-object type, evaluating an `attribute_reference` for the `Access` or `Unchecked_Access` attribute of an aliased view of an object. Non-null values of an access-to-subprogram type are obtained by evaluating an `attribute_reference` for the `Access` attribute of a non-intrinsic subprogram..
- 13.1/2 A `null_exclusion` in a construct specifies that the null value does not belong to the access subtype defined by the construct, that is, the access subtype *excludes null*. In addition, the anonymous access subtype defined by the `access_definition` for a controlling access parameter (see 3.9.2) excludes null. Finally, for a `subtype_indication` without a `null_exclusion`, the subtype denoted by the `subtype_indication` excludes null if and only if the subtype denoted by the `subtype_mark` in the `subtype_indication` excludes null.
- 14/1 All subtypes of an access-to-subprogram type are constrained. The first subtype of a type defined by an `access_definition` or an `access_to_object_definition` is unconstrained if the designated subtype is an unconstrained array or discriminated subtype; otherwise it is constrained.

#### Legality Rules

- 14.1/2 If a `subtype_indication`, `discriminant_specification`, `parameter_specification`, `parameter_and_result_profile`, `object_renaming_declaration`, or `formal_object_declaration` has a `null_exclusion`, the `subtype_mark` in that construct shall denote an access subtype that does not exclude null.

## Dynamic Semantics

A **composite\_constraint** is *compatible* with an unconstrained access subtype if it is compatible with the designated subtype. A **null\_exclusion** is compatible with any access subtype that does not exclude null. An access value *satisfies* a **composite\_constraint** of an access subtype if it equals the null value of its type or if it designates an object whose value satisfies the constraint. An access value satisfies an exclusion of the null value if it does not equal the null value of its type.

The elaboration of an **access\_type\_definition** creates the access type and its first subtype. For an access-to-object type, this elaboration includes the elaboration of the **subtype\_indication**, which creates the designated subtype.

The elaboration of an **access\_definition** creates an anonymous access type.

## NOTES

80 Access values are called “pointers” or “references” in some other languages.

81 Each access-to-object type has an associated storage pool; several access types can share the same pool. An object can be created in the storage pool of an access type by an allocator (see 4.8) for the access type. A storage pool (roughly) corresponds to what some other languages call a “heap.” See 13.11 for a discussion of pools.

82 Only **index\_constraints** and **discriminant\_constraints** can be applied to access types (see 3.6.1 and 3.7.1).

## Examples

*Examples of access-to-object types:*

```
type Peripheral_Ref is not null access Peripheral; -- see 3.8.1
type Binop_Ptr is access all Binary_Operation'Class;
-- general access-to-class-wide, see 3.9.1
```

*Example of an access subtype:*

```
subtype Drum_Ref is Peripheral_Ref(Drum); -- see 3.8.1
```

*Example of an access-to-subprogram type:*

```
type Message_Procedure is access procedure (M : in String := "Error!");
procedure Default_Message_Procedure(M : in String);
Give_Message : Message_Procedure := Default_Message_Procedure'Access;
...
procedure Other_Procedure(M : in String);
...
Give_Message := Other_Procedure'Access;
...
Give_Message("File not found."); -- call with parameter (.all is optional)
Give_Message.all; -- call with no parameters
```

### 3.10.1 Incomplete Type Declarations

There are no particular limitations on the designated type of an access type. In particular, the type of a component of the designated type can be another access type, or even the same access type. This permits mutually dependent and recursive access types. An **incomplete\_type\_declaration** can be used to introduce a type to be used as a designated type, while deferring its full definition to a subsequent **full\_type\_declaration**.

## Syntax

incomplete\_type\_declaration ::= **type** defining\_identifier [**discriminant\_part**] [**is tagged**];

*Static Semantics*

- 2.1/2 An *incomplete\_type\_declaration* declares an *incomplete view* of a type and its first subtype; the first subtype is unconstrained if a *discriminant\_part* appears. If the *incomplete\_type\_declaration* includes the reserved word **tagged**, it declares a *tagged incomplete view*. An incomplete view of a type is a limited view of the type (see 7.5).
- 2.2/2 Given an access type *A* whose designated type *T* is an incomplete view, a dereference of a value of type *A* also has this incomplete view except when:
- 2.3/2 • it occurs within the immediate scope of the completion of *T*, or
  - 2.4/2 • it occurs within the scope of a *nonlimited\_with\_clause* that mentions a library package in whose visible part the completion of *T* is declared.
- 2.5/2 In these cases, the dereference has the full view of *T*.
- 2.6/2 Similarly, if a *subtype\_mark* denotes a *subtype\_declaration* defining a subtype of an incomplete view *T*, the *subtype\_mark* denotes an incomplete view except under the same two circumstances given above, in which case it denotes the full view of *T*.

*Legality Rules*

- 3 An *incomplete\_type\_declaration* requires a completion, which shall be a *full\_type\_declaration*. If the *incomplete\_type\_declaration* occurs immediately within either the visible part of a *package\_specification* or a *declarative\_part*, then the *full\_type\_declaration* shall occur later and immediately within this visible part or *declarative\_part*. If the *incomplete\_type\_declaration* occurs immediately within the private part of a given *package\_specification*, then the *full\_type\_declaration* shall occur later and immediately within either the private part itself, or the *declarative\_part* of the corresponding *package\_body*.
- 4/2 If an *incomplete\_type\_declaration* includes the reserved word **tagged**, then a *full\_type\_declaration* that completes it shall declare a tagged type. If an *incomplete\_type\_declaration* has a *known\_discriminant\_part*, then a *full\_type\_declaration* that completes it shall have a fully conforming (explicit) *known\_discriminant\_part* (see 6.3.1). If an *incomplete\_type\_declaration* has no *discriminant\_part* (or an *unknown\_discriminant\_part*), then a corresponding *full\_type\_declaration* is nevertheless allowed to have discriminants, either explicitly, or inherited via derivation.
- 5/2 A name that denotes an incomplete view of a type may be used as follows:
- 6 • as the *subtype\_mark* in the *subtype\_indication* of an *access\_to\_object\_definition*; the only form of constraint allowed in this *subtype\_indication* is a *discriminant\_constraint*;
  - 7/2 • as the *subtype\_mark* in the *subtype\_indication* of a *subtype\_declaration*; the *subtype\_indication* shall not have a *null\_exclusion* or a *constraint*;
  - 8/2 • as the *subtype\_mark* in an *access\_definition*.
- 8.1/2 If such a name denotes a tagged incomplete view, it may also be used:
- 8.2/2 • as the *subtype\_mark* defining the subtype of a parameter in a *formal\_part*;
  - 9/2 • as the prefix of an *attribute\_reference* whose *attribute\_designator* is *Class*; such an *attribute\_reference* is restricted to the uses allowed here; it denotes a tagged incomplete view.
- 9.1/2 If such a name occurs within the declaration list containing the completion of the incomplete view, it may also be used:
- 9.2/2 • as the *subtype\_mark* defining the subtype of a parameter or result of an *access\_to\_subprogram\_definition*.

If any of the above uses occurs as part of the declaration of a primitive subprogram of the incomplete view, and the declaration occurs immediately within the private part of a package, then the completion of the incomplete view shall also occur immediately within the private part; it shall not be deferred to the package body.

No other uses of a name that denotes an incomplete view of a type are allowed.

9.3/2

A prefix that denotes an object shall not be of an incomplete view.

10/2

#### Static Semantics

*This paragraph was deleted.*

11/2

#### Dynamic Semantics

The elaboration of an `incomplete_type_declaration` has no effect.

12

#### NOTES

83 Within a `declarative_part`, an `incomplete_type_declaration` and a corresponding `full_type_declaration` cannot be separated by an intervening body. This is because a type has to be completely defined before it is frozen, and a body freezes all types declared prior to it in the same `declarative_part` (see 13.14).

13

#### Examples

*Example of a recursive type:*

14

```

type Cell; -- incomplete type declaration
type Link is access Cell;
type Cell is
 record
 Value : Integer;
 Succ : Link;
 Pred : Link;
 end record;
Head : Link := new Cell'(0, null, null);
Next : Link := Head.Succ;
```

15

16

17

*Examples of mutually dependent access types:*

18

```

type Person(<>); -- incomplete type declaration
type Car is tagged; -- incomplete type declaration
type Person_Name is access Person;
type Car_Name is access all Car'Class;
type Car is tagged
 record
 Number : Integer;
 Owner : Person_Name;
 end record;
type Person(Sex : Gender) is
 record
 Name : String(1 .. 20);
 Birth : Date;
 Age : Integer range 0 .. 130;
 Vehicle : Car_Name;
 case Sex is
 when M => Wife : Person_Name(Sex => F);
 when F => Husband : Person_Name(Sex => M);
 end case;
 end record;
My_Car, Your_Car, Next_Car : Car_Name := new Car; -- see 4.8
George : Person_Name := new Person(M);
...
George.Vehicle := Your_Car;
```

19/2

20/2

21/2

22

23

### 3.10.2 Operations of Access Types

- 1 The attribute `Access` is used to create access values designating aliased objects and non-intrinsic subprograms. The “accessibility” rules prevent dangling references (in the absence of uses of certain unchecked features — see Section 13).

*Name Resolution Rules*

- 2/2 For an `attribute_reference` with `attribute_designator Access` (or `Unchecked_Access` — see 13.10), the expected type shall be a single access type *A* such that:
- 2.1/2 • *A* is an access-to-object type with designated type *D* and the type of the prefix is *D'Class* or is covered by *D*, or
  - 2.2/2 • *A* is an access-to-subprogram type whose designated profile is type conformant with that of the prefix.
- 2.3/2 The prefix of such an `attribute_reference` is never interpreted as an `implicit_dereference` or a parameterless `function_call` (see 4.1.4). The designated type or profile of the expected type of the `attribute_reference` is the expected type or profile for the prefix.

*Static Semantics*

- 3/2 The accessibility rules, which prevent dangling references, are written in terms of *accessibility levels*, which reflect the run-time nesting of *masters*. As explained in 7.6.1, a master is the execution of a certain construct, such as a `subprogram_body`. An accessibility level is *deeper than* another if it is more deeply nested at run time. For example, an object declared local to a called subprogram has a deeper accessibility level than an object declared local to the calling subprogram. The accessibility rules for access types require that the accessibility level of an object designated by an access value be no deeper than that of the access type. This ensures that the object will live at least as long as the access type, which in turn ensures that the access value cannot later designate an object that no longer exists. The `Unchecked_Access` attribute may be used to circumvent the accessibility rules.
- 4 A given accessibility level is said to be *statically deeper* than another if the given level is known at compile time (as defined below) to be deeper than the other for all possible executions. In most cases, accessibility is enforced at compile time by Legality Rules. Run-time accessibility checks are also used, since the Legality Rules do not cover certain cases involving access parameters and generic packages.
- 5 Each master, and each entity and view created by it, has an accessibility level:
- 6 • The accessibility level of a given master is deeper than that of each dynamically enclosing master, and deeper than that of each master upon which the task executing the given master directly depends (see 9.3).
- 7/2 • An entity or view defined by a declaration and created as part of its elaboration has the same accessibility level as the innermost master of the declaration except in the cases of renaming and derived access types described below. A parameter of a master has the same accessibility level as the master.
- 8 • The accessibility level of a view of an object or subprogram defined by a `renaming_declaration` is the same as that of the renamed view.
- 9/2 • The accessibility level of a view conversion, `qualified_expression`, or parenthesized expression, is the same as that of the operand.
- 10/2 • The accessibility level of an `aggregate` or the result of a function call (or equivalent use of an operator) that is used (in its entirety) to directly initialize part of an object is that of the object

being initialized. In other contexts, the accessibility level of an aggregate or the result of a function call is that of the innermost master that evaluates the aggregate or function call.

- Within a return statement, the accessibility level of the return object is that of the execution of the return statement. If the return statement completes normally by returning from the function, then prior to leaving the function, the accessibility level of the return object changes to be a level determined by the point of call, as does the level of any coextensions (see below) of the return object. 10.1/2
  - The accessibility level of a derived access type is the same as that of its ultimate ancestor. 11
  - The accessibility level of the anonymous access type defined by an `access_definition` of an `object_renaming_declaration` is the same as that of the renamed view. 11.1/2
  - The accessibility level of the anonymous access type of an access discriminant in the `subtype_indication` or `qualified_expression` of an allocator, or in the `expression` or `return_subtype_indication` of a return statement is determined as follows: 12/2
    - If the value of the access discriminant is determined by a `discriminant_association` in a `subtype_indication`, the accessibility level of the object or subprogram designated by the associated value (or library level if the value is null); 12.1/2
    - If the value of the access discriminant is determined by a `record_component_association` in an aggregate, the accessibility level of the object or subprogram designated by the associated value (or library level if the value is null); 12.2/2
    - In other cases, where the value of the access discriminant is determined by an object with an unconstrained nominal subtype, the accessibility level of the object. 12.3/2
  - The accessibility level of the anonymous access type of an access discriminant in any other context is that of the enclosing object. 12.4/2
  - The accessibility level of the anonymous access type of an access parameter specifying an access-to-object type is the same as that of the view designated by the actual. 13/2
  - The accessibility level of the anonymous access type of an access parameter specifying an access-to-subprogram type is deeper than that of any master; all such anonymous access types have this same level. 13.1/2
  - The accessibility level of an object created by an allocator is the same as that of the access type, except for an allocator of an anonymous access type that defines the value of an access parameter or an access discriminant. For an allocator defining the value of an access parameter, the accessibility level is that of the innermost master of the call. For one defining an access discriminant, the accessibility level is determined as follows: 14/2
    - for an allocator used to define the constraint in a `subtype_declaration`, the level of the `subtype_declaration`; 14.1/2
    - for an allocator used to define the constraint in a `component_definition`, the level of the enclosing type; 14.2/2
    - for an allocator used to define the discriminant of an object, the level of the object. 14.3/2
- In this last case, the allocated object is said to be a *coextension* of the object whose discriminant designates it, as well as of any object of which the discriminated object is itself a coextension or subcomponent. All coextensions of an object are finalized when the object is finalized (see 7.6.1). 14.4/2
- The accessibility level of a view of an object or subprogram denoted by a dereference of an access value is the same as that of the access type. 15

- 16     • The accessibility level of a component, protected subprogram, or entry of (a view of) a composite object is the same as that of (the view of) the composite object.
- 16.1/2 In the above rules, the operand of a view conversion, parenthesized expression or `qualified_expression` is considered to be used in a context if the view conversion, parenthesized expression or `qualified_expression` itself is used in that context.
- 17 One accessibility level is defined to be *statically deeper* than another in the following cases:
- 18     • For a master that is statically nested within another master, the accessibility level of the inner master is statically deeper than that of the outer master.
  - 18.1/2     • The accessibility level of the anonymous access type of an access parameter specifying an access-to-subprogram type is statically deeper than that of any master; all such anonymous access types have this same level.
  - 19.2     • The statically deeper relationship does not apply to the accessibility level of the anonymous type of an access parameter specifying an access-to-object type; that is, such an accessibility level is not considered to be statically deeper, nor statically shallower, than any other.
  - 20     • For determining whether one level is statically deeper than another when within a generic package body, the generic package is presumed to be instantiated at the same level as where it was declared; run-time checks are needed in the case of more deeply nested instantiations.
  - 21     • For determining whether one level is statically deeper than another when within the declarative region of a `type_declaration`, the current instance of the type is presumed to be an object created at a deeper level than that of the type.
- 22 The accessibility level of all library units is called the *library level*; a library-level declaration or entity is one whose accessibility level is the library level.
- 23 The following attribute is defined for a prefix X that denotes an aliased view of an object:
- 24/1 `X'Access`     `X'Access` yields an access value that designates the object denoted by X. The type of `X'Access` is an access-to-object type, as determined by the expected type. The expected type shall be a general access type. X shall denote an aliased view of an object, including possibly the current instance (see 8.6) of a limited type within its definition, or a formal parameter or generic formal object of a tagged type. The view denoted by the prefix X shall satisfy the following additional requirements, presuming the expected type for `X'Access` is the general access type A with designated type D:
- 25     • If A is an access-to-variable type, then the view shall be a variable; on the other hand, if A is an access-to-constant type, the view may be either a constant or a variable.
  - 26/2     • The view shall not be a subcomponent that depends on discriminants of a variable whose nominal subtype is unconstrained, unless this subtype is indefinite, or the variable is constrained by its initial value.
  - 27/2     • If A is a named access type and D is a tagged type, then the type of the view shall be covered by D; if A is anonymous and D is tagged, then the type of the view shall be either D'Class or a type covered by D; if D is untagged, then the type of the view shall be D, and either:
    - 27.1/2         • the designated subtype of A shall statically match the nominal subtype of the view; or
    - 27.2/2         • D shall be discriminated in its full view and unconstrained in any partial view, and the designated subtype of A shall be unconstrained.

- The accessibility level of the view shall not be statically deeper than that of the access type *A*. In addition to the places where Legality Rules normally apply (see 12.3), this rule applies also in the private part of an instance of a generic unit. 28
- A check is made that the accessibility level of X is not deeper than that of the access type *A*. If this check fails, Program\_Error is raised. 29
- If the nominal subtype of X does not statically match the designated subtype of *A*, a view conversion of X to the designated subtype is evaluated (which might raise Constraint\_Error — see 4.6) and the value of X'Access designates that view. 30
- The following attribute is defined for a prefix P that denotes a subprogram: 31
- P'Access      P'Access yields an access value that designates the subprogram denoted by P. The type of P'Access is an access-to-subprogram type (*S*), as determined by the expected type. The accessibility level of P shall not be statically deeper than that of *S*. In addition to the places where Legality Rules normally apply (see 12.3), this rule applies also in the private part of an instance of a generic unit. The profile of P shall be subtype-conformant with the designated profile of *S*, and shall not be Intrinsic. If the subprogram denoted by P is declared within a generic unit, and the expression P'Access occurs within the body of that generic unit or within the body of a generic unit declared within the declarative region of the generic unit, then the ultimate ancestor of *S* shall be either a non-formal type declared within the generic unit or an anonymous access type of an access parameter. 32/2
- NOTES**
- 84 The Unchecked\_Access attribute yields the same result as the Access attribute for objects, but has fewer restrictions (see 13.10). There are other predefined operations that yield access values: an allocator can be used to create an object, and return an access value that designates it (see 4.8); evaluating the literal **null** yields a null access value that designates no entity at all (see 4.2). 33
- 85 The predefined operations of an access type also include the assignment operation, qualification, and membership tests. Explicit conversion is allowed between general access types with matching designated subtypes; explicit conversion is allowed between access-to-subprogram types with subtype conformant profiles (see 4.6). Named access types have predefined equality operators; anonymous access types do not, but they can use the predefined equality operators for *universal\_access* (see 4.5.2). 34/2
- 86 The object or subprogram designated by an access value can be named with a dereference, either an *explicit\_dereference* or an *implicit\_dereference*. See 4.1. 35
- 87 A call through the dereference of an access-to-subprogram value is never a dispatching call. 36
- 88 The Access attribute for subprograms and parameters of an anonymous access-to-subprogram type may together be used to implement “downward closures” — that is, to pass a more-nested subprogram as a parameter to a less-nested subprogram, as might be appropriate for an iterator abstraction or numerical integration. Downward closures can also be implemented using generic formal subprograms (see 12.6). Note that Unchecked\_Access is not allowed for subprograms. 37/2
- 89 Note that using an access-to-class-wide tagged type with a dispatching operation is a potentially more structured alternative to using an access-to-subprogram type. 38
- 90 An implementation may consider two access-to-subprogram values to be unequal, even though they designate the same subprogram. This might be because one points directly to the subprogram, while the other points to a special prologue that performs an Elaboration\_Check and then jumps to the subprogram. See 4.5.2. 39

*Examples**Example of use of the Access attribute:*

```
Martha : Person_Name := new Person(F);
Cars : array(1..2) of aliased Car;
...
Martha.Vehicle := Cars(1)'Access;
George.Vehicle := Cars(2)'Access;
```

## 3.11 Declarative Parts

- 1 A declarative\_part contains declarative\_items (possibly none).

*Syntax*

```

2 declarative_part ::= {declarative_item}
3 declarative_item ::=
4/1 basic_declarative_item | body
 basic_declarative_item ::=
 basic_declaration | aspect_clause | use_clause
5 body ::= proper_body | body_stub
6 proper_body ::=
 subprogram_body | package_body | task_body | protected_body

```

*Static Semantics*

- 6.1/2 The list of declarative\_items of a declarative\_part is called the *declaration list* of the declarative\_part.

*Dynamic Semantics*

- 7 The elaboration of a declarative\_part consists of the elaboration of the declarative\_items, if any, in the order in which they are given in the declarative\_part.
- 8 An elaborable construct is in the *elaborated* state after the normal completion of its elaboration. Prior to that, it is *not yet elaborated*.
- 9 For a construct that attempts to use a body, a check (Elaboration\_Check) is performed, as follows:
- 10/1 • For a call to a (non-protected) subprogram that has an explicit body, a check is made that the body is already elaborated. This check and the evaluations of any actual parameters of the call are done in an arbitrary order.
- 11 • For a call to a protected operation of a protected type (that has a body — no check is performed if a pragma Import applies to the protected type), a check is made that the protected\_body is already elaborated. This check and the evaluations of any actual parameters of the call are done in an arbitrary order.
- 12 • For the activation of a task, a check is made by the activator that the task\_body is already elaborated. If two or more tasks are being activated together (see 9.2), as the result of the elaboration of a declarative\_part or the initialization for the object created by an allocator, this check is done for all of them before activating any of them.
- 13 • For the instantiation of a generic unit that has a body, a check is made that this body is already elaborated. This check and the evaluation of any explicit\_generic\_actual\_parameters of the instantiation are done in an arbitrary order.
- 14 The exception Program\_Error is raised if any of these checks fails.

### 3.11.1 Completions of Declarations

- 1/1 Declarations sometimes come in two parts. A declaration that requires a second part is said to *require completion*. The second part is called the *completion* of the declaration (and of the entity declared), and is either another declaration, a body, or a pragma. A *body* is a *body*, an *entry\_body*, or a renaming-as-body (see 8.5.4).

*Name Resolution Rules*

- A construct that can be a completion is interpreted as the completion of a prior declaration only if:
- The declaration and the completion occur immediately within the same declarative region;
  - The defining name or `defining_program_unit_name` in the completion is the same as in the declaration, or in the case of a `pragma`, the `pragma` applies to the declaration;
  - If the declaration is overloadable, then the completion either has a type-conformant profile, or is a `pragma`.

*Legality Rules*

An implicit declaration shall not have a completion. For any explicit declaration that is specified to *require completion*, there shall be a corresponding explicit completion.

At most one completion is allowed for a given declaration. Additional requirements on completions appear where each kind of completion is defined.

A type is *completely defined* at a place that is after its full type definition (if it has one) and after all of its subcomponent types are completely defined. A type shall be completely defined before it is frozen (see 13.14 and 7.3).

## NOTES

91 Completions are in principle allowed for any kind of explicit declaration. However, for some kinds of declaration, the only allowed completion is a `pragma Import`, and implementations are not required to support `pragma Import` for every kind of entity.

92 There are rules that prevent premature uses of declarations that have a corresponding completion. The Elaboration\_Checks of 3.11 prevent such uses at run time for subprograms, protected operations, tasks, and generic units. The rules of 13.14, “Freezing Rules” prevent, at compile time, premature uses of other entities such as private types and deferred constants.

## Section 4: Names and Expressions

The rules applicable to the different forms of name and expression, and to their evaluation, are given in this section.

### 4.1 Names

Names can denote declared entities, whether declared explicitly or implicitly (see 3.1). Names can also denote objects or subprograms designated by access values; the results of type\_conversions or function\_calls; subcomponents and slices of objects and values; protected subprograms, single entries, entry families, and entries in families of entries. Finally, names can denote attributes of any of the foregoing.

| <i>Syntax</i>                                             | 1 |
|-----------------------------------------------------------|---|
| <code>name ::=</code>                                     |   |
| <code>direct_name             explicit_dereference</code> | 2 |
| <code>  indexed_component     slice</code>                |   |
| <code>  selected_component   attribute_reference</code>   |   |
| <code>  type_conversion      function_call</code>         |   |
| <code>  character_literal</code>                          |   |
| <code>direct_name ::= identifier   operator_symbol</code> | 3 |
| <code>prefix ::= name   implicit_dereference</code>       | 4 |
| <code>explicit_dereference ::= name.all</code>            | 5 |
| <code>implicit_dereference ::= name</code>                | 6 |

Certain forms of name (indexed\_components, selected\_components, slices, and attribute\_references) include a prefix that is either itself a name that denotes some related entity, or an implicit\_dereference of an access value that designates some related entity.

#### *Name Resolution Rules*

The name in a dereference (either an implicit\_dereference or an explicit\_dereference) is expected to be of any access type.

#### *Static Semantics*

If the type of the name in a dereference is some access-to-object type  $T$ , then the dereference denotes a view of an object, the nominal subtype of the view being the designated subtype of  $T$ .

If the type of the name in a dereference is some access-to-subprogram type  $S$ , then the dereference denotes a view of a subprogram, the profile of the view being the designated profile of  $S$ .

#### *Dynamic Semantics*

The evaluation of a name determines the entity denoted by the name. This evaluation has no other effect for a name that is a direct\_name or a character\_literal.

The evaluation of a name that has a prefix includes the evaluation of the prefix. The evaluation of a prefix consists of the evaluation of the name or the implicit\_dereference. The prefix denotes the entity denoted by the name or the implicit\_dereference.

- 13 The evaluation of a dereference consists of the evaluation of the name and the determination of the object or subprogram that is designated by the value of the name. A check is made that the value of the name is not the null access value. Constraint\_Error is raised if this check fails. The dereference denotes the object or subprogram designated by the value of the name.

*Examples*

- 14 Examples of direct names:

|        |                                         |             |
|--------|-----------------------------------------|-------------|
| Pi     | -- the direct name of a number          | (see 3.3.2) |
| Limit  | -- the direct name of a constant        | (see 3.3.1) |
| Count  | -- the direct name of a scalar variable | (see 3.3.1) |
| Board  | -- the direct name of an array variable | (see 3.6.1) |
| Matrix | -- the direct name of a type            | (see 3.6)   |
| Random | -- the direct name of a function        | (see 6.1)   |
| Error  | -- the direct name of an exception      | (see 11.1)  |

- 16 Examples of dereferences:

|                |                                                                                                           |
|----------------|-----------------------------------------------------------------------------------------------------------|
| Next_Car.all   | -- explicit dereference denoting the object designated by<br>-- the access variable Next_Car (see 3.10.1) |
| Next_Car.Owner | -- selected component with implicit dereference;<br>-- same as Next_Car.all.Owner                         |

### 4.1.1 Indexed Components

- 1 An indexed\_component denotes either a component of an array or an entry in a family of entries.

*Syntax*

2 indexed\_component ::= prefix(expression {, expression})

*Name Resolution Rules*

- 3 The prefix of an indexed\_component with a given number of expressions shall resolve to denote an array (after any implicit dereference) with the corresponding number of index positions, or shall resolve to denote an entry family of a task or protected object (in which case there shall be only one expression).
- 4 The expected type for each expression is the corresponding index type.

*Static Semantics*

- 5 When the prefix denotes an array, the indexed\_component denotes the component of the array with the specified index value(s). The nominal subtype of the indexed\_component is the component subtype of the array type.
- 6 When the prefix denotes an entry family, the indexed\_component denotes the individual entry of the entry family with the specified index value.

*Dynamic Semantics*

- 7 For the evaluation of an indexed\_component, the prefix and the expressions are evaluated in an arbitrary order. The value of each expression is converted to the corresponding index type. A check is made that each index value belongs to the corresponding index range of the array or entry family denoted by the prefix. Constraint\_Error is raised if this check fails.

| <i>Examples</i>                        | 8                                                    |
|----------------------------------------|------------------------------------------------------|
| <i>Examples of indexed components:</i> | 9                                                    |
| My_Schedule(Sat)                       | — a component of a one-dimensional array (see 3.6.1) |
| Page(10)                               | — a component of a one-dimensional array (see 3.6)   |
| Board(M, J + 1)                        | — a component of a two-dimensional array (see 3.6.1) |
| Page(10)(20)                           | — a component of a component (see 3.6)               |
| Request(Medium)                        | — an entry in a family of entries (see 9.1)          |
| Next_Frame(L)(M, N)                    | — a component of a function call (see 6.1)           |

## NOTES

1 *Notes on the examples:* Distinct notations are used for components of multidimensional arrays (such as Board) and arrays of arrays (such as Page). The components of an array of arrays are arrays and can therefore be indexed. Thus Page(10)(20) denotes the 20th component of Page(10). In the last example Next\_Frame(L) is a function call returning an access value that designates a two-dimensional array.

10

**4.1.2 Slices**

A **slice** denotes a one-dimensional array formed by a sequence of consecutive components of a one-dimensional array. A **slice** of a variable is a variable; a **slice** of a constant is a constant; a **slice** of a value is a value.

1

*Syntax*

slice ::= prefix(discrete\_range)

2

*Name Resolution Rules*

The **prefix** of a **slice** shall resolve to denote a one-dimensional array (after any implicit dereference).

3

The expected type for the **discrete\_range** of a **slice** is the index type of the array type.

4

*Static Semantics*

A **slice** denotes a one-dimensional array formed by the sequence of consecutive components of the array denoted by the **prefix**, corresponding to the range of values of the index given by the **discrete\_range**.

5

The type of the **slice** is that of the **prefix**. Its bounds are those defined by the **discrete\_range**.

6

*Dynamic Semantics*

For the evaluation of a **slice**, the **prefix** and the **discrete\_range** are evaluated in an arbitrary order. If the **slice** is not a *null slice* (a slice where the **discrete\_range** is a null range), then a check is made that the bounds of the **discrete\_range** belong to the index range of the array denoted by the **prefix**. **Constraint\_Error** is raised if this check fails.

7

## NOTES

2 A **slice** is not permitted as the **prefix** of an **Access attribute\_reference**, even if the components or the array as a whole are aliased. See 3.10.2.

8

3 For a one-dimensional array A, the **slice** A(N .. N) denotes an array that has only one component; its type is the type of A. On the other hand, A(N) denotes a component of the array A and has the corresponding component type.

9

*Examples*

10 Examples of slices:

|                       |                                   |                       |
|-----------------------|-----------------------------------|-----------------------|
| 11 Stars(1 .. 15)     | -- a slice of 15 characters       | (see 3.6.3)           |
| Page(10 .. 10 + Size) | -- a slice of 1 + Size components | (see 3.6)             |
| Page(L)(A .. B)       | -- a slice of the array Page(L)   | (see 3.6)             |
| Stars(1 .. 0)         | -- a null slice                   | (see 3.6.3)           |
| My_Schedule(Weekday)  | -- bounds given by subtype        | (see 3.6.1 and 3.5.1) |
| Stars(5 .. 15)(K)     | -- same as Stars(K)               | (see 3.6.3)           |
|                       | -- provided that K is in 5 .. 15  |                       |

### 4.1.3 Selected Components

1 Selected\_components are used to denote components (including discriminants), entries, entry families, and protected subprograms; they are also used as expanded names as described below.

*Syntax*2 selected\_component ::= prefix . selector\_name  
3 selector\_name ::= identifier | character\_literal | operator\_symbol*Name Resolution Rules*4 A selected\_component is called an *expanded name* if, according to the visibility rules, at least one possible interpretation of its prefix denotes a package or an enclosing named construct (directly, not through a subprogram\_renaming\_declaration or generic\_renaming\_declaration).

5 A selected\_component that is not an expanded name shall resolve to denote one of the following:

- 6 • A component (including a discriminant):

7 The prefix shall resolve to denote an object or value of some non-array composite type (after any implicit dereference). The selector\_name shall resolve to denote a discriminant\_specification of the type, or, unless the type is a protected type, a component\_declaration of the type. The selected\_component denotes the corresponding component of the object or value.

- 8 • A single entry, an entry family, or a protected subprogram:

9 The prefix shall resolve to denote an object or value of some task or protected type (after any implicit dereference). The selector\_name shall resolve to denote an entry\_declaration or subprogram\_declaration occurring (implicitly or explicitly) within the visible part of that type. The selected\_component denotes the corresponding entry, entry family, or protected subprogram.

- 9.1/2 • A view of a subprogram whose first formal parameter is of a tagged type or is an access parameter whose designated type is tagged:

9.2/2 The prefix (after any implicit dereference) shall resolve to denote an object or value of a specific tagged type *T* or class-wide type *TClass*. The selector\_name shall resolve to denote a view of a subprogram declared immediately within the declarative region in which an ancestor of the type *T* is declared. The first formal parameter of the subprogram shall be of type *T*, or a class-wide type that covers *T*, or an access parameter designating one of these types. The designator of the subprogram shall not be the same as that of a component of the tagged type visible at the point of the selected\_component. The selected\_component denotes a view of this subprogram that omits the first formal parameter. This view is called a *prefixed view* of the subprogram, and the prefix of the selected\_component (after any implicit dereference) is called the *prefix* of the prefixed view.

10 An expanded name shall resolve to denote a declaration that occurs immediately within a named declarative region, as follows:

- The prefix shall resolve to denote either a package (including the current instance of a generic package, or a rename of a package), or an enclosing named construct. 11
- The selector\_name shall resolve to denote a declaration that occurs immediately within the declarative region of the package or enclosing construct (the declaration shall be visible at the place of the expanded name — see 8.3). The expanded name denotes that declaration. 12
- If the prefix does not denote a package, then it shall be a direct\_name or an expanded name, and it shall resolve to denote a program unit (other than a package), the current instance of a type, a block\_statement, a loop\_statement, or an accept\_statement (in the case of an accept\_statement or entry\_body, no family index is allowed); the expanded name shall occur within the declarative region of this construct. Further, if this construct is a callable construct and the prefix denotes more than one such enclosing callable construct, then the expanded name is ambiguous, independently of the selector\_name. 13

#### *Legality Rules*

For a subprogram whose first parameter is an access parameter, the prefix of any prefixed view shall 13.1/2 denote an aliased view of an object.

For a subprogram whose first parameter is of mode **in** **out** or **out**, or of an anonymous access-to-variable 13.2/2 type, the prefix of any prefixed view shall denote a variable.

#### *Dynamic Semantics*

The evaluation of a selected\_component includes the evaluation of the prefix. 14

For a selected\_component that denotes a component of a variant, a check is made that the values of the discriminants are such that the value or object denoted by the prefix has this component. The exception Constraint\_Error is raised if this check fails. 15

#### *Examples*

##### *Examples of selected components:*

|                    |                                                         |              |
|--------------------|---------------------------------------------------------|--------------|
| Tomorrow.Month     | -- a record component                                   | (see 3.8)    |
| Next_Car.Owner     | -- a record component                                   | (see 3.10.1) |
| Next_Car.Owner.Age | -- a record component                                   | (see 3.10.1) |
|                    | -- the previous two lines involve implicit dereferences |              |
| Writer.Unit        | -- a record component (a discriminant)                  | (see 3.8.1)  |
| Min_Cell(H).Value  | -- a record component of the result                     | (see 6.1)    |
|                    | -- of the function call Min_Cell(H)                     |              |
| Cashier.Append     | -- a prefixed view of a procedure                       | (see 3.9.4)  |
| Control.Seize      | -- an entry of a protected object                       | (see 9.4)    |
| Pool(K).Write      | -- an entry of the task Pool(K)                         | (see 9.4)    |

##### *Examples of expanded names:*

|                  |                                                 |             |
|------------------|-------------------------------------------------|-------------|
| Key_Manager."<"  | -- an operator of the visible part of a package | (see 7.3.1) |
| Dot_Product.Sum  | -- a variable declared in a function body       | (see 6.1)   |
| Buffer.Pool      | -- a variable declared in a protected unit      | (see 9.11)  |
| Buffer.Read      | -- an entry of a protected unit                 | (see 9.11)  |
| Swap.Temp        | -- a variable declared in a block statement     | (see 5.6)   |
| Standard.Boolean | -- the name of a predefined type                | (see A.1)   |

## 4.1.4 Attributes

- 1 An *attribute* is a characteristic of an entity that can be queried via an *attribute\_reference* or a *range\_attribute\_reference*.

### Syntax

- 2 *attribute\_reference* ::= prefix'attribute\_designator
- 3   *attribute\_designator* ::=  
    identifier[*static\_expression*]  
    | Access | Delta | Digits
- 4   *range\_attribute\_reference* ::= prefix'range\_attribute\_designator
- 5   *range\_attribute\_designator* ::= Range[*static\_expression*]

### Name Resolution Rules

- 6 In an *attribute\_reference*, if the *attribute\_designator* is for an attribute defined for (at least some) objects of an access type, then the prefix is never interpreted as an *implicit\_dereference*; otherwise (and for all *range\_attribute\_references*), if the type of the name within the prefix is of an access type, the prefix is interpreted as an *implicit\_dereference*. Similarly, if the *attribute\_designator* is for an attribute defined for (at least some) functions, then the prefix is never interpreted as a parameterless *function\_call*; otherwise (and for all *range\_attribute\_references*), if the prefix consists of a name that denotes a function, it is interpreted as a parameterless *function\_call*.
- 7 The expression, if any, in an *attribute\_designator* or *range\_attribute\_designator* is expected to be of any integer type.

### Legality Rules

- 8 The expression, if any, in an *attribute\_designator* or *range\_attribute\_designator* shall be static.

### Static Semantics

- 9 An *attribute\_reference* denotes a value, an object, a subprogram, or some other kind of program entity.
- 10 A *range\_attribute\_reference* X'Range(N) is equivalent to the *range* X'First(N) .. X'Last(N), except that the prefix is only evaluated once. Similarly, X'Range is equivalent to X'First .. X'Last, except that the prefix is only evaluated once.

### Dynamic Semantics

- 11 The evaluation of an *attribute\_reference* (or *range\_attribute\_reference*) consists of the evaluation of the prefix.

### Implementation Permissions

- 12/1 An implementation may provide implementation-defined attributes; the identifier for an implementation-defined attribute shall differ from those of the language-defined attributes unless supplied for compatibility with a previous edition of this International Standard.

#### NOTES

- 13 4 Attributes are defined throughout this International Standard, and are summarized in Annex K.
- 14/2 5 In general, the name in a prefix of an *attribute\_reference* (or a *range\_attribute\_reference*) has to be resolved without using any context. However, in the case of the Access attribute, the expected type for the *attribute\_reference* has to be a single access type, and the resolution of the name can use the fact that the type of the object or the profile of the callable

entity denoted by the prefix has to match the designated type or be type conformant with the designated profile of the access type.

#### Examples

##### *Examples of attributes:*

|                    |                                                 |             |    |
|--------------------|-------------------------------------------------|-------------|----|
| Color'First        | -- minimum value of the enumeration type Color  | (see 3.5.1) | 15 |
| Rainbow'Base'First | -- same as Color'First                          | (see 3.5.1) | 16 |
| Real'Digits        | -- precision of the type Real                   | (see 3.5.7) |    |
| Board'Last(2)      | -- upper bound of the second dimension of Board | (see 3.6.1) |    |
| Board'Range(1)     | -- index range of the first dimension of Board  | (see 3.6.1) |    |
| Pool(K)'Terminated | -- True if task Pool(K) is terminated           | (see 9.1)   |    |
| Date'Size          | -- number of bits for records of type Date      | (see 3.8)   |    |
| Message'Address    | -- address of the record variable Message       | (see 3.7.1) |    |

## 4.2 Literals

A *literal* represents a value literally, that is, by means of notation suited to its kind. A literal is either a *numeric\_literal*, a *character\_literal*, the literal **null**, or a *string\_literal*.

#### Name Resolution Rules

*This paragraph was deleted.*

For a name that consists of a *character\_literal*, either its expected type shall be a single character type, in which case it is interpreted as a parameterless *function\_call* that yields the corresponding value of the character type, or its expected profile shall correspond to a parameterless function with a character result type, in which case it is interpreted as the name of the corresponding parameterless function declared as part of the character type's definition (see 3.5.1). In either case, the *character\_literal* denotes the *enumeration\_literal\_specification*.

The expected type for a primary that is a *string\_literal* shall be a single string type.

#### Legality Rules

A *character\_literal* that is a name shall correspond to a *defining\_character\_literal* of the expected type, or of the result type of the expected profile.

For each character of a *string\_literal* with a given expected string type, there shall be a corresponding *defining\_character\_literal* of the component type of the expected string type.

*This paragraph was deleted.*

#### Static Semantics

An integer literal is of type *universal\_integer*. A real literal is of type *universal\_real*. The literal **null** is of type *universal\_access*.

#### Dynamic Semantics

The evaluation of a numeric literal, or the literal **null**, yields the represented value.

The evaluation of a *string\_literal* that is a primary yields an array value containing the value of each character of the sequence of characters of the *string\_literal*, as defined in 2.6. The bounds of this array value are determined according to the rules for *positional\_array\_aggregates* (see 4.3.3), except that for a null string literal, the upper bound is the predecessor of the lower bound.

- 11 For the evaluation of a `string_literal` of type `T`, a check is made that the value of each character of the `string_literal` belongs to the component subtype of `T`. For the evaluation of a null string literal, a check is made that its lower bound is greater than the lower bound of the base range of the index type. The exception `Constraint_Error` is raised if either of these checks fails.

12 **NOTES**

6 Enumeration literals that are identifiers rather than `character_literals` follow the normal rules for identifiers when used in a name (see 4.1 and 4.1.3). `Character_literals` used as `selector_names` follow the normal rules for expanded names (see 4.1.3).

*Examples*

- 13 Examples of literals:

14    3.14159\_26536    -- a real literal  
     1\_345            -- an integer literal  
     'À'              -- a character literal  
     "Some Text"     -- a string literal

## 4.3 Aggregates

- 1 An `aggregate` combines component values into a composite value of an array type, record type, or record extension.

*Syntax*

2 `aggregate ::= record_aggregate | extension_aggregate | array_aggregate`

*Name Resolution Rules*

- 3/2 The expected type for an `aggregate` shall be a single array type, record type, or record extension.

*Legality Rules*

- 4 An `aggregate` shall not be of a class-wide type.

*Dynamic Semantics*

- 5 For the evaluation of an `aggregate`, an anonymous object is created and values for the components or ancestor part are obtained (as described in the subsequent subclause for each kind of the `aggregate`) and assigned into the corresponding components or ancestor part of the anonymous object. Obtaining the values and the assignments occur in an arbitrary order. The value of the `aggregate` is the value of this object.
- 6 If an `aggregate` is of a tagged type, a check is made that its value belongs to the first subtype of the type. `Constraint_Error` is raised if this check fails.

### 4.3.1 Record Aggregates

- 1 In a `record_aggregate`, a value is specified for each component of the record or record extension value, using either a named or a positional association.

*Syntax*

2 `record_aggregate ::= (record_component_association_list)`  
3    `record_component_association_list ::=`  
      `record_component_association {, record_component_association}`  
      `| null record`

```

record_component_association ::= 4/2
 [component_choice_list =>] expression
 | component_choice_list => <>
component_choice_list ::= 5
 component_selector_name { | component_selector_name}
 | others

```

A record\_component\_association is a *named component association* if it has a component\_choice\_list; otherwise, it is a *positional component association*. Any positional component associations shall precede any named component associations. If there is a named association with a component\_choice\_list of others, it shall come last.

In the record\_component\_association\_list for a record\_aggregate, if there is only one association, it shall be a named association.

#### *Name Resolution Rules*

The expected type for a record\_aggregate shall be a single record type or record extension.

For the record\_component\_association\_list of a record\_aggregate, all components of the composite value defined by the aggregate are *needed*; for the association list of an extension\_aggregate, only those components not determined by the ancestor expression or subtype are needed (see 4.3.2). Each selector\_name in a record\_component\_association shall denote a needed component (including possibly a discriminant).

The expected type for the expression of a record\_component\_association is the type of the *associated* component(s); the associated component(s) are as follows:

- For a positional association, the component (including possibly a discriminant) in the corresponding relative position (in the declarative region of the type), counting only the needed components;
- For a named association with one or more component\_selector\_names, the named component(s);
- For a named association with the reserved word others, all needed components that are not associated with some previous association.

#### *Legality Rules*

If the type of a record\_aggregate is a record extension, then it shall be a descendant of a record type, through one or more record extensions (and no private extensions).

If there are no components needed in a given record\_component\_association\_list, then the reserved words **null record** shall appear rather than a list of record\_component\_associations.

Each record\_component\_association other than an others choice with a <> shall have at least one associated component, and each needed component shall be associated with exactly one record\_component\_association. If a record\_component\_association with an expression has two or more associated components, all of them shall be of the same type.

If the components of a variant\_part are needed, then the value of a discriminant that governs the variant\_part shall be given by a static expression.

A record\_component\_association for a discriminant without a default\_expression shall have an expression rather than <>.

Dynamic Semantics

- 18 The evaluation of a `record_aggregate` consists of the evaluation of the `record_component_association_list`.
  - 19 For the evaluation of a `record_component_association_list`, any per-object constraints (see 3.8) for components specified in the association list are elaborated and any `expressions` are evaluated and converted to the subtype of the associated component. Any constraint elaborations and `expression` evaluations (and conversions) occur in an arbitrary order, except that the `expression` for a discriminant is evaluated (and converted) prior to the elaboration of any per-object constraint that depends on it, which in turn occurs prior to the evaluation and conversion of the `expression` for the component with the per-object constraint.
  - 19.1/2 For a `record_component_association` with an `expression`, the `expression` defines the value for the associated component(s). For a `record_component_association` with `<>`, if the `component_declaration` has a `default_expression`, that `default_expression` defines the value for the associated component(s); otherwise, the associated component(s) are initialized by default as for a stand-alone object of the component subtype (see 3.3.1).
  - 20 The `expression` of a `record_component_association` is evaluated (and converted) once for each associated component.

## NOTES

- <sup>21</sup> 7 For a record\_aggregate with positional associations, expressions specifying discriminant values appear first since the known\_discriminant\_part is given first in the declaration of the type; they have to be in the same order as in the known discriminant part.

### *Examples*

- ```

22 Example of a record aggregate with positional associations:
23   (4, July, 1776) -- see 3.8

24 Examples of record aggregates with named associations:
25   (Day => 4, Month => July, Year => 1776)
      (Month => July, Day => 4, Year => 1776)

26   (Disk, Closed, Track => 5, Cylinder => 12) -- see 3.8.1
      (Unit => Disk, Status => Closed, Cylinder => 9, Track => 1)

27/2 Examples of component associations with several choices:
28   (Value => 0, Succ|Pred => new Cell'(0, null, null)) -- see 3.10.1
29     -- The allocator is evaluated twice: Succ and Pred designate different cells
29.1/2 (Value => 0, Succ|Pred => <>) -- see 3.10.1
29.2/2   -- Succ and Pred will be set to null

30 Examples of record aggregates for tagged types (see 3.9 and 3.9.1):
31   Expression'(null record)
   Literal'(Value => 0.0)
   Painted Point'(0.0, Pi/2.0, Paint => Red)

```

4.3.2 Extension Aggregates

An extension_aggregate specifies a value for a type that is a record extension by specifying a value or subtype for an ancestor of the type, followed by associations for any components not determined by the ancestor_part.

Syntax

```
extension_aggregate ::= 1
  (ancestor_part with record_component_association_list)
  ancestor_part ::= expression | subtype_mark 2
```

Name Resolution Rules

The expected type for an extension_aggregate shall be a single type that is a record extension. If the ancestor_part is an expression, it is expected to be of any tagged type.

Legality Rules

If the ancestor_part is a subtype_mark, it shall denote a specific tagged subtype. If the ancestor_part is an expression, it shall not be dynamically tagged. The type of the extension_aggregate shall be derived from the type of the ancestor_part, through one or more record extensions (and no private extensions).

Static Semantics

For the record_component_association_list of an extension_aggregate, the only components *needed* are those of the composite value defined by the aggregate that are not inherited from the type of the ancestor_part, plus any inherited discriminants if the ancestor_part is a subtype_mark that denotes an unconstrained subtype.

Dynamic Semantics

For the evaluation of an extension_aggregate, the record_component_association_list is evaluated. If the ancestor_part is an expression, it is also evaluated; if the ancestor_part is a subtype_mark, the components of the value of the aggregate not given by the record_component_association_list are initialized by default as for an object of the ancestor type. Any implicit initializations or evaluations are performed in an arbitrary order, except that the expression for a discriminant is evaluated prior to any other evaluation or initialization that depends on it.

If the type of the ancestor_part has discriminants that are not inherited by the type of the extension_aggregate, then, unless the ancestor_part is a subtype_mark that denotes an unconstrained subtype, a check is made that each discriminant of the ancestor has the value specified for a corresponding discriminant, either in the record_component_association_list, or in the derived_type_definition for some ancestor of the type of the extension_aggregate. Constraint_Error is raised if this check fails.

NOTES

8 If all components of the value of the extension_aggregate are determined by the ancestor_part, then the record_component_association_list is required to be simply null record.

9 If the ancestor_part is a subtype_mark, then its type can be abstract. If its type is controlled, then as the last step of evaluating the aggregate, the Initialize procedure of the ancestor type is called, unless the Initialize procedure is abstract (see 7.6).

Examples

- 11 Examples of extension aggregates (for types defined in 3.9.1):
- 12 Painted_Point'(Point **with** Red)
 (Point'(*P*) **with** Paint => Black)
- 13 (Expression **with** Left => 1.2, Right => 3.4)
 Addition'(Binop **with** null record)
 -- presuming Binop is of type Binary_Operation

4.3.3 Array Aggregates

- 1 In an array_aggregate, a value is specified for each component of an array, either positionally or by its index. For a positional_array_aggregate, the components are given in increasing-index order, with a final **others**, if any, representing any remaining components. For a named_array_aggregate, the components are identified by the values covered by the discrete_choices.

Syntax

- 2 array_aggregate ::=
 positional_array_aggregate | named_array_aggregate
- 3/2 positional_array_aggregate ::=
 (expression, expression {, expression})
 | (expression {, expression}, others => expression)
 | (expression {, expression}, others => <>)
- 4 named_array_aggregate ::=
 (array_component_association {, array_component_association})
- 5/2 array_component_association ::=
 discrete_choice_list => expression
 | discrete_choice_list => <>

- 6 An *n-dimensional array_aggregate* is one that is written as n levels of nested array_aggregates (or at the bottom level, equivalent string_literals). For the multidimensional case ($n \geq 2$) the array_aggregates (or equivalent string_literals) at the $n-1$ lower levels are called *subaggregates* of the enclosing *n-dimensional array_aggregate*. The expressions of the bottom level subaggregates (or of the array_aggregate itself if one-dimensional) are called the *array component expressions* of the enclosing *n-dimensional array_aggregate*.

Name Resolution Rules

- 7/2 The expected type for an array_aggregate (that is not a subaggregate) shall be a single array type. The component type of this array type is the expected type for each array component expression of the array_aggregate.

- 8 The expected type for each discrete_choice in any discrete_choice_list of a named_array_aggregate is the type of the *corresponding index*; the corresponding index for an array_aggregate that is not a subaggregate is the first index of its type; for an $(n-m)$ -dimensional subaggregate within an array_aggregate of an *n-dimensional type*, the corresponding index is the index in position $m+1$.

Legality Rules

- 9 An array_aggregate of an *n-dimensional array type* shall be written as an *n-dimensional array_aggregate*.
- 10 An **others** choice is allowed for an array_aggregate only if an *applicable index constraint* applies to the array_aggregate. An applicable index constraint is a constraint provided by certain contexts where an

`array_aggregate` is permitted that can be used to determine the bounds of the array value specified by the aggregate. Each of the following contexts (and none other) defines an applicable index constraint:

- For an `explicit_actual_parameter`, an `explicit_generic_actual_parameter`, the expression of a return statement, the initialization expression in an `object_declaration`, or a `default_expression` (for a parameter or a component), when the nominal subtype of the corresponding formal parameter, generic formal parameter, function return object, object, or component is a constrained array subtype, the applicable index constraint is the constraint of the subtype; 11/2
- For the expression of an `assignment_statement` where the name denotes an array variable, the applicable index constraint is the constraint of the array variable; 12
- For the operand of a `qualified_expression` whose `subtype_mark` denotes a constrained array subtype, the applicable index constraint is the constraint of the subtype; 13
- For a component expression in an `aggregate`, if the component's nominal subtype is a constrained array subtype, the applicable index constraint is the constraint of the subtype; 14
- For a parenthesized expression, the applicable index constraint is that, if any, defined for the expression. 15

The applicable index constraint *applies* to an `array_aggregate` that appears in such a context, as well as to any subaggregates thereof. In the case of an `explicit_actual_parameter` (or `default_expression`) for a call on a generic formal subprogram, no applicable index constraint is defined. 16

The `discrete_choice_list` of an `array_component_association` is allowed to have a `discrete_choice` that is a nonstatic expression or that is a `discrete_range` that defines a nonstatic or null range, only if it is the single `discrete_choice` of its `discrete_choice_list`, and there is only one `array_component_association` in the `array_aggregate`. 17

In a `named_array_aggregate` with more than one `discrete_choice`, no two `discrete_choices` are allowed to cover the same value (see 3.8.1); if there is no `others` choice, the `discrete_choices` taken together shall exactly cover a contiguous sequence of values of the corresponding index type. 18

A bottom level subaggregate of a multidimensional `array_aggregate` of a given array type is allowed to be a `string_literal` only if the component type of the array type is a character type; each character of such a `string_literal` shall correspond to a `defining_character_literal` of the component type. 19

Static Semantics

A subaggregate that is a `string_literal` is equivalent to one that is a `positional_array_aggregate` of the same length, with each expression being the `character_literal` for the corresponding character of the `string_literal`. 20

Dynamic Semantics

The evaluation of an `array_aggregate` of a given array type proceeds in two steps: 21

1. Any `discrete_choices` of this aggregate and of its subaggregates are evaluated in an arbitrary order, and converted to the corresponding index type; 22
2. The array component expressions of the aggregate are evaluated in an arbitrary order and their values are converted to the component subtype of the array type; an array component expression is evaluated once for each associated component. 23

Each expression in an `array_component_association` defines the value for the associated component(s). For an `array_component_association` with `<>`, the associated component(s) are initialized by default as for a stand-alone object of the component subtype (see 3.3.1). 23.1/2

- 24 The bounds of the index range of an `array_aggregate` (including a subaggregate) are determined as follows:
- 25 • For an `array_aggregate` with an `others` choice, the bounds are those of the corresponding index range from the applicable index constraint;
 - 26 • For a `positional_array_aggregate` (or equivalent `string_literal`) without an `others` choice, the lower bound is that of the corresponding index range in the applicable index constraint, if defined, or that of the corresponding index subtype, if not; in either case, the upper bound is determined from the lower bound and the number of `expressions` (or the length of the `string_literal`);
 - 27 • For a `named_array_aggregate` without an `others` choice, the bounds are determined by the smallest and largest index values covered by any `discrete_choice_list`.
- 28 For an `array_aggregate`, a check is made that the index range defined by its bounds is compatible with the corresponding index subtype.
- 29 For an `array_aggregate` with an `others` choice, a check is made that no `expression` is specified for an index value outside the bounds determined by the applicable index constraint.
- 30 For a multidimensional `array_aggregate`, a check is made that all subaggregates that correspond to the same index have the same bounds.
- 31 The exception `Constraint_Error` is raised if any of the above checks fail.

NOTES

32/2 10 In an `array_aggregate`, positional notation may only be used with two or more `expressions`; a single `expression` in parentheses is interpreted as a parenthesized expression. A `named_array_aggregate`, such as `(1 => X)`, may be used to specify an array with a single component.

Examples

33 Examples of array aggregates with positional associations:

34 `(7, 9, 5, 1, 3, 2, 4, 8, 6, 0)`
`Table'(5, 8, 4, 1, others => 0)` -- see 3.6

35 Examples of array aggregates with named associations:

36 `(1 .. 5 => (1 .. 8 => 0.0))` -- two-dimensional
`(1 .. N => new Cell)` -- *N* new cells, in particular for *N* = 0

37 `Table'(2 | 4 | 10 => 1, others => 0)`
`Schedule'(Mon .. Fri => True, others => False)` -- see 3.6
`Schedule'(Wed | Sun => False, others => True)`
`Vector'(1 => 2.5)` -- single-component vector

38 Examples of two-dimensional array aggregates:

39 -- Three aggregates for the same value of subtype `Matrix(1..2,1..3)` (see 3.6):
`((1.1, 1.2, 1.3), (2.1, 2.2, 2.3))`
`(1 => (1.1, 1.2, 1.3), 2 => (2.1, 2.2, 2.3))`
`(1 => (1 => 1.1, 2 => 1.2, 3 => 1.3), 2 => (1 => 2.1, 2 => 2.2, 3 => 2.3))`

41 Examples of aggregates as initial values:

42 `A : Table := (7, 9, 5, 1, 3, 2, 4, 8, 6, 0);` -- *A(1)=7, A(10)=0*
`B : Table := (2 | 4 | 10 => 1, others => 0);` -- *B(1)=0, B(10)=1*
`C : constant Matrix := (1 .. 5 => (1 .. 8 => 0.0));` -- *C'Last(1)=5, C'Last(2)=8*

43 `D : Bit_Vector(M .. N) := (M .. N => True);` -- see 3.6
`E : Bit_Vector(M .. N) := (others => True);`
`F : String(1 .. 1) := (1 => 'F');` -- a one component aggregate: same as "F"

Example of an array aggregate with defaulted others choice and with an applicable index constraint provided by an enclosing record aggregate: 44/2

Buffer'(Size => 50, Pos => 1, Value => String'('x', others => <>)) -- see 3.7 45/2

4.4 Expressions

An *expression* is a formula that defines the computation or retrieval of a value. In this International Standard, the term “expression” refers to a construct of the syntactic category **expression** or of any of the other five syntactic categories defined below. 1

Syntax

```

expression ::= 2
  relation {and relation} | relation {and then relation}
  | relation {or relation} | relation {or else relation}
  | relation {xor relation}

relation ::= 3
  simple_expression [relational_operator simple_expression]
  | simple_expression [not] in range
  | simple_expression [not] in subtype_mark

simple_expression ::= [unary_adding_operator] term {binary_adding_operator term} 4

term ::= factor {multiplying_operator factor} 5

factor ::= primary [** primary] | abs primary | not primary 6

primary ::= 7
  numeric_literal | null | string_literal | aggregate
  | name | qualified_expression | allocator | (expression)

```

Name Resolution Rules

A **name** used as a **primary** shall resolve to denote an object or a value. 8

Static Semantics

Each expression has a type; it specifies the computation or retrieval of a value of that type. 9

Dynamic Semantics

The value of a **primary** that is a **name** denoting an object is the value of the object. 10

Implementation Permissions

For the evaluation of a **primary** that is a **name** denoting an object of an unconstrained numeric subtype, if the value of the object is outside the base range of its type, the implementation may either raise Constraint_Error or return the value of the object. 11

Examples

12 Examples of primaries:

13	4.0	-- real literal
	Pi	-- named number
	(1 .. 10 => 0)	-- array aggregate
	Sum	-- variable
	Integer'Last	-- attribute
	Sine(X)	-- function call
	Color'(Blue)	-- qualified expression
	Real(M*N)	-- conversion
	(Line_Count + 10)	-- parenthesized expression

14 Examples of expressions:

15/2	Volume	-- primary
	not Destroyed	-- factor
	2*Line_Count	-- term
	-4.0	-- simple expression
	-4.0 + A	-- simple expression
	B**2 - 4.0*A*C	-- simple expression
	R*Sin(θ)*Cos(φ)	-- simple expression
	Password(1 .. 3) = "Bvv"	-- relation
	Count in Small_Int	-- relation
	Count not in Small_Int	-- relation
	Index = 0 or Item_Hit	-- expression
	(Cold and Sunny) or Warm	-- expression (parentheses are required)
	A**(B**C)	-- expression (parentheses are required)

4.5 Operators and Expression Evaluation

- 1 The language defines the following six categories of operators (given in order of increasing precedence). The corresponding operator_symbols, and only those, can be used as designators in declarations of functions for user-defined operators. See 6.6, “Overloading of Operators”.

Syntax

2	logical_operator ::=	and or xor
3	relational_operator ::=	= /= < <= > >=
4	binary_adding_operator ::=	+ - &
5	unary_adding_operator ::=	+ -
6	multiplying_operator ::=	* / mod rem
7	highest_precedence_operator ::=	** abs not

Static Semantics

- 8 For a sequence of operators of the same precedence level, the operators are associated with their operands in textual order from left to right. Parentheses can be used to impose specific associations.
- 9 For each form of type definition, certain of the above operators are *predefined*; that is, they are implicitly declared immediately after the type definition. For each such implicit operator declaration, the parameters are called Left and Right for *binary* operators; the single parameter is called Right for *unary* operators. An expression of the form X op Y, where op is a binary operator, is equivalent to a function_call of the form "op"(X, Y). An expression of the form op Y, where op is a unary operator, is equivalent to a function_call of the form "op"(Y). The predefined operators and their effects are described in subclauses 4.5.1 through 4.5.6.

Dynamic Semantics

The predefined operations on integer types either yield the mathematically correct result or raise the exception `Constraint_Error`. For implementations that support the Numerics Annex, the predefined operations on real types yield results whose accuracy is defined in Annex G, or raise the exception `Constraint_Error`.

Implementation Requirements

The implementation of a predefined operator that delivers a result of an integer or fixed point type may raise `Constraint_Error` only if the result is outside the base range of the result type.

The implementation of a predefined operator that delivers a result of a floating point type may raise `Constraint_Error` only if the result is outside the safe range of the result type.

Implementation Permissions

For a sequence of predefined operators of the same precedence level (and in the absence of parentheses imposing a specific association), an implementation may impose any association of the operators with operands so long as the result produced is an allowed result for the left-to-right association, but ignoring the potential for failure of language-defined checks in either the left-to-right or chosen order of association.

NOTES

11 The two operands of an expression of the form `X op Y`, where `op` is a binary operator, are evaluated in an arbitrary order, as for any `function_call` (see 6.4).

*Examples**Examples of precedence:*

<code>not Sunny or Warm</code>	-- same as <i>(not Sunny) or Warm</i>	15
<code>X > 4.0 and Y > 0.0</code>	-- same as <i>(X > 4.0) and (Y > 0.0)</i>	16
<code>-4.0*A**2</code>	-- same as <i>-(4.0 * (A**2))</i>	17
<code>abs(1 + A) + B</code>	-- same as <i>(abs (1 + A)) + B</i>	
<code>Y**(-3)</code>	-- parentheses are necessary	
<code>A / B * C</code>	-- same as <i>(A/B)*C</i>	
<code>A + (B + C)</code>	-- evaluate <i>B + C</i> before adding it to <i>A</i>	

4.5.1 Logical Operators and Short-circuit Control Forms

Name Resolution Rules

An expression consisting of two relations connected by `and` `then` or `or else` (a *short-circuit control form*) shall resolve to be of some boolean type; the expected type for both relations is that same boolean type.

Static Semantics

The following logical operators are predefined for every boolean type *T*, for every modular type *T*, and for every one-dimensional array type *T* whose component type is a boolean type:

```
function "and" (Left, Right : T) return T
function "or" (Left, Right : T) return T
function "xor" (Left, Right : T) return T
```

For boolean types, the predefined logical operators `and`, `or`, and `xor` perform the conventional operations of conjunction, inclusive disjunction, and exclusive disjunction, respectively.

For modular types, the predefined logical operators are defined on a bit-by-bit basis, using the binary representation of the value of the operands to yield a binary representation for the result, where zero

represents False and one represents True. If this result is outside the base range of the type, a final subtraction by the modulus is performed to bring the result into the base range of the type.

- 6 The logical operators on arrays are performed on a component-by-component basis on matching components (as for equality — see 4.5.2), using the predefined logical operator for the component type. The bounds of the resulting array are those of the left operand.

Dynamic Semantics

- 7 The short-circuit control forms **and then** and **or else** deliver the same result as the corresponding predefined **and** and **or** operators for boolean types, except that the left operand is always evaluated first, and the right operand is not evaluated if the value of the left operand determines the result.
- 8 For the logical operators on arrays, a check is made that for each component of the left operand there is a matching component of the right operand, and vice versa. Also, a check is made that each component of the result belongs to the component subtype. The exception **Constraint_Error** is raised if either of the above checks fails.

NOTES

- 9 12 The conventional meaning of the logical operators is given by the following truth table:

A	B	(A and B)	(A or B)	(A xor B)
True	True	True	True	False
True	False	False	True	True
False	True	False	True	True
False	False	False	False	False

Examples

- 11 Examples of logical operators:

```
Sunny or Warm
Filter(1 .. 10) and Filter(15 .. 24)    -- see 3.6.1
```

- 13 Examples of short-circuit control forms:

```
Next_Car.Owner /= null and then Next_Car.Owner.Age > 25    -- see 3.10.1
N = 0 or else A(N) = Hit_Value
```

4.5.2 Relational Operators and Membership Tests

- 1 The *equality operators* = (equals) and /= (not equals) are predefined for nonlimited types. The other *relational_operators* are the *ordering operators* < (less than), <= (less than or equal), > (greater than), and >= (greater than or equal). The ordering operators are predefined for scalar types, and for *discrete_array_types*, that is, one-dimensional array types whose components are of a discrete type.
- 2 A *membership test*, using **in** or **not in**, determines whether or not a value belongs to a given subtype or range, or has a tag that identifies a type that is covered by a given type. Membership tests are allowed for all types.

Name Resolution Rules

- 3/2 The *tested type* of a membership test is the type of the **range** or the type determined by the **subtype_mark**. If the tested type is tagged, then the **simple_expression** shall resolve to be of a type that is convertible (see 4.6) to the tested type; if untagged, the expected type for the **simple_expression** is the tested type.

Legality Rules

For a membership test, if the `simple_expression` is of a tagged class-wide type, then the tested type shall
be (visibly) tagged.

Static Semantics

The result type of a membership test is the predefined type Boolean.

The equality operators are predefined for every specific type T that is not limited, and not an anonymous access type, with the following specifications:

```
function "=" (Left, Right : T) return Boolean
function "/=" (Left, Right : T) return Boolean
```

The following additional equality operators for the `universal_access` type are declared in package Standard
for use with anonymous access types:

```
function "=" (Left, Right : universal_access) return Boolean
function "/=" (Left, Right : universal_access) return Boolean
```

The ordering operators are predefined for every specific scalar type T , and for every discrete array type T ,
with the following specifications:

```
function "<" (Left, Right : T) return Boolean
function "<=" (Left, Right : T) return Boolean
function ">" (Left, Right : T) return Boolean
function ">=" (Left, Right : T) return Boolean
```

Name Resolution Rules

At least one of the operands of an equality operator for `universal_access` shall be of a specific anonymous access type. Unless the predefined equality operator is identified using an expanded name with prefix denoting the package Standard, neither operand shall be of an access-to-object type whose designated type is D or D 'Class, where D has a user-defined primitive equality operator such that:

- its result type is Boolean;
- it is declared immediately within the same declaration list as D ; and
- at least one of its operands is an access parameter with designated type D .

Legality Rules

At least one of the operands of the equality operators for `universal_access` shall be of type `universal_access`, or both shall be of access-to-object types, or both shall be of access-to-subprogram types. Further:

- When both are of access-to-object types, the designated types shall be the same or one shall cover the other, and if the designated types are elementary or array types, then the designated subtypes shall statically match;
- When both are of access-to-subprogram types, the designated profiles shall be subtype conformant.

Dynamic Semantics

For discrete types, the predefined relational operators are defined in terms of corresponding mathematical operations on the position numbers of the values of the operands.

For real types, the predefined relational operators are defined in terms of the corresponding mathematical operations on the values of the operands, subject to the accuracy of the type.

- 12 Two access-to-object values are equal if they designate the same object, or if both are equal to the null value of the access type.
- 13 Two access-to-subprogram values are equal if they are the result of the same evaluation of an Access attribute_reference, or if both are equal to the null value of the access type. Two access-to-subprogram values are unequal if they designate different subprograms. It is unspecified whether two access values that designate the same subprogram but are the result of distinct evaluations of Access attribute_references are equal or unequal.
- 14 For a type extension, predefined equality is defined in terms of the primitive (possibly user-defined) equals operator of the parent type and of any tagged components of the extension part, and predefined equality for any other components not inherited from the parent type.
- 15 For a private type, if its full type is tagged, predefined equality is defined in terms of the primitive equals operator of the full type; if the full type is untagged, predefined equality for the private type is that of its full type.
- 16 For other composite types, the predefined equality operators (and certain other predefined operations on composite types — see 4.5.1 and 4.6) are defined in terms of the corresponding operation on *matching components*, defined as follows:
- 17 • For two composite objects or values of the same non-array type, matching components are those that correspond to the same component_declaration or discriminant_specification;
 - 18 • For two one-dimensional arrays of the same type, matching components are those (if any) whose index values match in the following sense: the lower bounds of the index ranges are defined to match, and the successors of matching indices are defined to match;
 - 19 • For two multidimensional arrays of the same type, matching components are those whose index values match in successive index positions.
- 20 The analogous definitions apply if the types of the two objects or values are convertible, rather than being the same.
- 21 Given the above definition of matching components, the result of the predefined equals operator for composite types (other than for those composite types covered earlier) is defined as follows:
- 22 • If there are no components, the result is defined to be True;
 - 23 • If there are unmatched components, the result is defined to be False;
 - 24 • Otherwise, the result is defined in terms of the primitive equals operator for any matching tagged components, and the predefined equals for any matching untagged components.
- 24.1/1 For any composite type, the order in which "`=`" is called for components is unspecified. Furthermore, if the result can be determined before calling "`=`" on some components, it is unspecified whether "`=`" is called on those components.
- 25 The predefined "`/=`" operator gives the complementary result to the predefined "`=`" operator.
- 26 For a discrete array type, the predefined ordering operators correspond to *lexicographic order* using the predefined order relation of the component type: A null array is lexicographically less than any array having at least one component. In the case of nonnull arrays, the left operand is lexicographically less than the right operand if the first component of the left operand is less than that of the right; otherwise the left operand is lexicographically less than the right operand only if their first components are equal and the tail of the left operand is lexicographically less than that of the right (the *tail* consists of the remaining components beyond the first and can be null).

For the evaluation of a membership test, the `simple_expression` and the `range` (if any) are evaluated in an arbitrary order. 27

A membership test using `in` yields the result True if: 28

- The tested type is scalar, and the value of the `simple_expression` belongs to the given `range`, or 29
the range of the named subtype; or
- The tested type is not scalar, and the value of the `simple_expression` satisfies any constraints of 30/2
the named subtype, and:
 - if the type of the `simple_expression` is class-wide, the value has a tag that identifies a type 30.1/2
covered by the tested type;
 - if the tested type is an access type and the named subtype excludes null, the value of the 30.2/2
`simple_expression` is not null.

Otherwise the test yields the result False. 31

A membership test using `not in` gives the complementary result to the corresponding membership test 32
using `in`.

Implementation Requirements

For all nonlimited types declared in language-defined packages, the "`=`" and "`/=`" operators of the type 32.1/1
shall behave as if they were the predefined equality operators for the purposes of the equality of composite
types and generic formal types.

NOTES

This paragraph was deleted.

13 If a composite type has components that depend on discriminants, two values of this type have matching components 34
if and only if their discriminants are equal. Two nonnull arrays have matching components if and only if the length of
each dimension is the same for both.

Examples

Examples of expressions involving relational operators and membership tests: 35

<code>X /= Y</code>	36
<code>" " < "A" and "A" < "Aa"</code>	37
<code>"Aa" < "B" and "A" < "A "</code>	37
<code>My_Car = null</code>	38
<code>My_Car = Your_Car</code>	38
<code>My_Car.all = Your_Car.all</code>	38
<code>N not in 1 .. 10</code>	39
<code>Today in Mon .. Fri</code>	39
<code>Today in Weekday</code>	39
<code>Archive in Disk_Unit</code>	39
<code>Tree.all in Addition'Class</code>	39
<code>-- true</code>	
<code>-- true</code>	
<code>-- true if My_Car has been set to null (see 3.10.1)</code>	
<code>-- true if we both share the same car</code>	
<code>-- true if the two cars are identical</code>	
<code>-- range membership test</code>	
<code>-- range membership test</code>	
<code>-- subtype membership test (see 3.5.1)</code>	
<code>-- subtype membership test (see 3.8.1)</code>	
<code>-- class membership test (see 3.9.1)</code>	

4.5.3 Binary Adding Operators

Static Semantics

The binary adding operators `+` (addition) and `-` (subtraction) are predefined for every specific numeric 1
type `T` with their conventional meaning. They have the following specifications:

```
function "+"(Left, Right : T) return T
function "-"(Left, Right : T) return T
```

- 3 The concatenation operators `&` are predefined for every nonlimited, one-dimensional array type T with component type C . They have the following specifications:

```
4   function "&"(Left : T; Right : T) return T
  function "&"(Left : T; Right : C) return T
  function "&"(Left : C; Right : T) return T
  function "&"(Left : C; Right : C) return T
```

Dynamic Semantics

- 5 For the evaluation of a concatenation with result type T , if both operands are of type T , the result of the concatenation is a one-dimensional array whose length is the sum of the lengths of its operands, and whose components comprise the components of the left operand followed by the components of the right operand. If the left operand is a null array, the result of the concatenation is the right operand. Otherwise, the lower bound of the result is determined as follows:

- 6 • If the ultimate ancestor of the array type was defined by a `constrained_array_definition`, then the lower bound of the result is that of the index subtype;
- 7 • If the ultimate ancestor of the array type was defined by an `unconstrained_array_definition`, then the lower bound of the result is that of the left operand.

- 8 The upper bound is determined by the lower bound and the length. A check is made that the upper bound of the result of the concatenation belongs to the range of the index subtype, unless the result is a null array. `Constraint_Error` is raised if this check fails.

- 9 If either operand is of the component type C , the result of the concatenation is given by the above rules, using in place of such an operand an array having this operand as its only component (converted to the component subtype) and having the lower bound of the index subtype of the array type as its lower bound.

- 10 The result of a concatenation is defined in terms of an assignment to an anonymous object, as for any function call (see 6.5).

NOTES

- 11 14 As for all predefined operators on modular types, the binary adding operators `+` and `-` on modular types include a final reduction modulo the modulus if the result is outside the base range of the type.

Examples

- 12 Examples of expressions involving binary adding operators:

```
13   Z + 0.1      -- Z has to be of a real type
14   "A" & "BCD"  -- concatenation of two string literals
   'A' & "BCD"   -- concatenation of a character literal and a string literal
   'A' & 'A'      -- concatenation of two character literals
```

4.5.4 Unary Adding Operators

Static Semantics

- 1 The unary adding operators `+` (identity) and `-` (negation) are predefined for every specific numeric type T with their conventional meaning. They have the following specifications:

```
2   function "+"(Right : T) return T
  function "-"(Right : T) return T
```

NOTES

- 3 15 For modular integer types, the unary adding operator `-`, when given a nonzero operand, returns the result of subtracting the value of the operand from the modulus; for a zero operand, the result is zero.

4.5.5 Multiplying Operators

Static Semantics

The multiplying operators `*` (multiplication), `/` (division), `mod` (modulus), and `rem` (remainder) are predefined for every specific integer type T :

```
function "*" (Left, Right : T) return T
function "/" (Left, Right : T) return T
function "mod"(Left, Right : T) return T
function "rem"(Left, Right : T) return T
```

Signed integer multiplication has its conventional meaning.

Signed integer division and remainder are defined by the relation:

$$A = (A/B) * B + (A \text{ rem } B)$$

where $(A \text{ rem } B)$ has the sign of A and an absolute value less than the absolute value of B . Signed integer division satisfies the identity:

$$(-A)/B = - (A/B) = A/(-B)$$

The signed integer modulus operator is defined such that the result of $A \text{ mod } B$ has the sign of B and an absolute value less than the absolute value of B ; in addition, for some signed integer value N , this result satisfies the relation:

$$A = B*N + (A \text{ mod } B)$$

The multiplying operators on modular types are defined in terms of the corresponding signed integer operators, followed by a reduction modulo the modulus if the result is outside the base range of the type (which is only possible for the `"*"` operator).

Multiplication and division operators are predefined for every specific floating point type T :

```
function "*" (Left, Right : T) return T
function "/" (Left, Right : T) return T
```

The following multiplication and division operators, with an operand of the predefined type `Integer`, are predefined for every specific fixed point type T :

```
function "*" (Left : T; Right : Integer) return T
function "*" (Left : Integer; Right : T) return T
function "/" (Left : T; Right : Integer) return T
```

All of the above multiplying operators are usable with an operand of an appropriate universal numeric type. The following additional multiplying operators for `root_real` are predefined, and are usable when both operands are of an appropriate universal or root numeric type, and the result is allowed to be of type `root_real`, as in a `number_declaration`:

```
function "*" (Left, Right : root_real) return root_real
function "/" (Left, Right : root_real) return root_real

function "*" (Left : root_real; Right : root_integer) return root_real
function "*" (Left : root_integer; Right : root_real) return root_real
function "/" (Left : root_real; Right : root_integer) return root_real
```

Multiplication and division between any two fixed point types are provided by the following two predefined operators:

```
function "*" (Left, Right : universal_fixed) return universal_fixed
function "/" (Left, Right : universal_fixed) return universal_fixed
```

Name Resolution Rules

- 19.1/2 The above two fixed-fixed multiplying operators shall not be used in a context where the expected type for the result is itself *universal_fixed* — the context has to identify some other numeric type to which the result is to be converted, either explicitly or implicitly. Unless the predefined universal operator is identified using an expanded name with prefix denoting the package Standard, an explicit conversion is required on the result when using the above fixed-fixed multiplication operator if either operand is of a type having a user-defined primitive multiplication operator such that:
- 19.2/2 • it is declared immediately within the same declaration list as the type; and
- 19.3/2 • both of its formal parameters are of a fixed-point type.
- 19.4/2 A corresponding requirement applies to the universal fixed-fixed division operator.

Legality Rules

- 20/2 *This paragraph was deleted.*

Dynamic Semantics

- 21 The multiplication and division operators for real types have their conventional meaning. For floating point types, the accuracy of the result is determined by the precision of the result type. For decimal fixed point types, the result is truncated toward zero if the mathematical result is between two multiples of the *small* of the specific result type (possibly determined by context); for ordinary fixed point types, if the mathematical result is between two multiples of the *small*, it is unspecified which of the two is the result.
- 22 The exception Constraint_Error is raised by integer division, **rem**, and **mod** if the right operand is zero. Similarly, for a real type T with $T'Machine_Overflows$ True, division by zero raises Constraint_Error.

NOTES

- 23 16 For positive A and B, A/B is the quotient and $A \bmod B$ is the remainder when A is divided by B. The following relations are satisfied by the rem operator:

$$\begin{aligned} A \bmod (-B) &= A \bmod B \\ (-A) \bmod B &= - (A \bmod B) \end{aligned}$$

- 25 17 For any signed integer K, the following identity holds:

$$A \bmod B = (A + K*B) \bmod B$$

27 The relations between signed integer division, remainder, and modulus are illustrated by the following table:

	A	B	A/B	$A \bmod B$	$A \bmod B$	A	B	A/B	$A \bmod B$	$A \bmod B$
29	10	5	2	0	0	-10	5	-2	0	0
	11	5	2	1	1	-11	5	-2	-1	4
	12	5	2	2	2	-12	5	-2	-2	3
	13	5	2	3	3	-13	5	-2	-3	2
	14	5	2	4	4	-14	5	-2	-4	1
30	A	B	A/B	$A \bmod B$	$A \bmod B$	A	B	A/B	$A \bmod B$	$A \bmod B$
	10	-5	-2	0	0	-10	-5	2	0	0
	11	-5	-2	1	-4	-11	-5	2	-1	-1
	12	-5	-2	2	-3	-12	-5	2	-2	-2
	13	-5	-2	3	-2	-13	-5	2	-3	-3
	14	-5	-2	4	-1	-14	-5	2	-4	-4

Examples

- 31 Examples of expressions involving multiplying operators:

```
I : Integer := 1;
J : Integer := 2;
K : Integer := 3;
```

X : Real := 1.0;	-- see 3.5.7	33
Y : Real := 2.0;		
F : Fraction := 0.25;	-- see 3.5.9	34
G : Fraction := 0.5;		
<i>Expression</i>	<i>Value</i>	<i>Result Type</i>
I*J	2	same as I and J, that is, Integer
K/J	1	same as K and J, that is, Integer
K mod J	1	same as K and J, that is, Integer
X/Y	0.5	same as X and Y, that is, Real
F/2	0.125	same as F, that is, Fraction
3*F	0.75	same as F, that is, Fraction
0.75*G	0.375	universal_fixed, implicitly convertible to any fixed point type
Fraction(F*G)	0.125	Fraction, as stated by the conversion
Real(J)*Y	4.0	Real, the type of both operands after conversion of J

4.5.6 Highest Precedence Operators

Static Semantics

The highest precedence unary operator **abs** (absolute value) is predefined for every specific numeric type *T*, with the following specification:

```
function "abs"(Right : T) return T
```

The highest precedence unary operator **not** (logical negation) is predefined for every boolean type *T*, every modular type *T*, and for every one-dimensional array type *T* whose components are of a boolean type, with the following specification:

```
function "not"(Right : T) return T
```

The result of the operator **not** for a modular type is defined as the difference between the high bound of the base range of the type and the value of the operand. For a binary modulus, this corresponds to a bit-wise complement of the binary representation of the value of the operand.

The operator **not** that applies to a one-dimensional array of boolean components yields a one-dimensional boolean array with the same bounds; each component of the result is obtained by logical negation of the corresponding component of the operand (that is, the component that has the same index value). A check is made that each component of the result belongs to the component subtype; the exception *Constraint_Error* is raised if this check fails.

The highest precedence *exponentiation* operator ****** is predefined for every specific integer type *T* with the following specification:

```
function "***"(Left : T; Right : Natural) return T
```

Exponentiation is also predefined for every specific floating point type as well as *root_real*, with the following specification (where *T* is *root_real* or the floating point type):

```
function "***"(Left : T; Right : Integer'Base) return T
```

The right operand of an exponentiation is the *exponent*. The expression *X**N* with the value of the exponent *N* positive is equivalent to the expression *X*X*...X* (with *N*-1 multiplications) except that the multiplications are associated in an arbitrary order. With *N* equal to zero, the result is one. With the value of *N* negative (only defined for a floating point operand), the result is the reciprocal of the result using the absolute value of *N* as the exponent.

Implementation Permissions

- 12 The implementation of exponentiation for the case of a negative exponent is allowed to raise `Constraint_Error` if the intermediate result of the repeated multiplications is outside the safe range of the type, even though the final result (after taking the reciprocal) would not be. (The best machine approximation to the final result in this case would generally be 0.0.)

NOTES

- 13 As implied by the specification given above for exponentiation of an integer type, a check is made that the exponent is not negative. `Constraint_Error` is raised if this check fails.

4.6 Type Conversions

- 1 Explicit type conversions, both value conversions and view conversions, are allowed between closely related types as defined below. This clause also defines rules for value and view conversions to a particular subtype of a type, both explicit ones and those implicit in other constructs.

Syntax

2 `type_conversion ::=`
 `subtype_mark(expression)`
 `| subtype_mark(name)`

- 3 The *target subtype* of a `type_conversion` is the subtype denoted by the `subtype_mark`. The *operand* of a `type_conversion` is the expression or name within the parentheses; its type is the *operand type*.

- 4 One type is *convertible* to a second type if a `type_conversion` with the first type as operand type and the second type as target type is legal according to the rules of this clause. Two types are convertible if each is convertible to the other.

- 5/2 A `type_conversion` whose operand is the *name* of an object is called a *view conversion* if both its target type and operand type are tagged, or if it appears in a call as an actual parameter of mode `out` or `in out`; other `type_conversions` are called *value conversions*.

Name Resolution Rules

- 6 The operand of a `type_conversion` is expected to be of any type.

- 7 The operand of a view conversion is interpreted only as a *name*; the operand of a value conversion is interpreted as an *expression*.

Legality Rules

- 8/2 In a view conversion for an untagged type, the target type shall be convertible (back) to the operand type.

Paragraphs 9 through 20 were reorganized and moved below.

- 21/2 If there is a type that is an ancestor of both the target type and the operand type, or both types are class-wide types, then at least one of the following rules shall apply:

- 21.1/2 • The target type shall be untagged; or
- 22 • The operand type shall be covered by or descended from the target type; or
- 23/2 • The operand type shall be a class-wide type that covers the target type; or
- 23.1/2 • The operand and target types shall both be class-wide types and the specific type associated with at least one of them shall be an interface type.

If there is no type that is the ancestor of both the target type and the operand type, and they are not both class-wide types, one of the following rules shall apply: 24/2

- If the target type is a numeric type, then the operand type shall be a numeric type. 24.1/2
- If the target type is an array type, then the operand type shall be an array type. Further:
 - The types shall have the same dimensionality; 24.3/2
 - Corresponding index types shall be convertible; 24.4/2
 - The component subtypes shall statically match; 24.5/2
 - If the component types are anonymous access types, then the accessibility level of the operand type shall not be statically deeper than that of the target type; 24.6/2
 - Neither the target type nor the operand type shall be limited; 24.7/2
 - If the target type of a view conversion has aliased components, then so shall the operand type; and 24.8/2
 - The operand type of a view conversion shall not have a tagged, private, or volatile subcomponent. 24.9/2
- If the target type is *universal_access*, then the operand type shall be an access type. 24.10/2
- If the target type is a general access-to-object type, then the operand type shall be *universal_access* or an access-to-object type. Further, if the operand type is not *universal_access*:
 - If the target type is an access-to-variable type, then the operand type shall be an access-to-variable type; 24.12/2
 - If the target designated type is tagged, then the operand designated type shall be convertible to the target designated type; 24.13/2
 - If the target designated type is not tagged, then the designated types shall be the same, and either:
 - the designated subtypes shall statically match; or 24.15/2
 - the designated type shall be discriminated in its full view and unconstrained in any partial view, and one of the designated subtypes shall be unconstrained; 24.16/2
 - The accessibility level of the operand type shall not be statically deeper than that of the target type. In addition to the places where Legality Rules normally apply (see 12.3), this rule applies also in the private part of an instance of a generic unit. 24.17/2
- If the target type is a pool-specific access-to-object type, then the operand type shall be *universal_access*. 24.18/2
- If the target type is an access-to-subprogram type, then the operand type shall be *universal_access* or an access-to-subprogram type. Further, if the operand type is not *universal_access*:
 - The designated profiles shall be subtype-conformant. 24.20/2
 - The accessibility level of the operand type shall not be statically deeper than that of the target type. In addition to the places where Legality Rules normally apply (see 12.3), this rule applies also in the private part of an instance of a generic unit. If the operand type is declared within a generic body, the target type shall be declared within the generic body. 24.21/2

Static Semantics

A *type_conversion* that is a value conversion denotes the value that is the result of converting the value of the operand to the target subtype. 25

- 26 A type_conversion that is a view conversion denotes a view of the object denoted by the operand. This view is a variable of the target type if the operand denotes a variable; otherwise it is a constant of the target type.
- 27 The nominal subtype of a type_conversion is its target subtype.

Dynamic Semantics

- 28 For the evaluation of a type_conversion that is a value conversion, the operand is evaluated, and then the value of the operand is converted to a corresponding value of the target type, if any. If there is no value of the target type that corresponds to the operand value, Constraint_Error is raised; this can only happen on conversion to a modular type, and only when the operand value is outside the base range of the modular type. Additional rules follow:
- 29 • Numeric Type Conversion
 - 30 • If the target and the operand types are both integer types, then the result is the value of the target type that corresponds to the same mathematical integer as the operand.
 - 31 • If the target type is a decimal fixed point type, then the result is truncated (toward 0) if the value of the operand is not a multiple of the *small* of the target type.
 - 32 • If the target type is some other real type, then the result is within the accuracy of the target type (see G.2, “Numeric Performance Requirements”, for implementations that support the Numerics Annex).
 - 33 • If the target type is an integer type and the operand type is real, the result is rounded to the nearest integer (away from zero if exactly halfway between two integers).
 - 34 • Enumeration Type Conversion
 - 35 • The result is the value of the target type with the same position number as that of the operand value.
 - 36 • Array Type Conversion
 - 37 • If the target subtype is a constrained array subtype, then a check is made that the length of each dimension of the value of the operand equals the length of the corresponding dimension of the target subtype. The bounds of the result are those of the target subtype.
 - 38 • If the target subtype is an unconstrained array subtype, then the bounds of the result are obtained by converting each bound of the value of the operand to the corresponding index type of the target type. For each nonnull index range, a check is made that the bounds of the range belong to the corresponding index subtype.
 - 39 • In either array case, the value of each component of the result is that of the matching component of the operand value (see 4.5.2).
 - 39.1/2 • If the component types of the array types are anonymous access types, then a check is made that the accessibility level of the operand type is not deeper than that of the target type.
 - 40 • Composite (Non-Array) Type Conversion
 - 41 • The value of each nondiscriminant component of the result is that of the matching component of the operand value.
 - 42 • The tag of the result is that of the operand. If the operand type is class-wide, a check is made that the tag of the operand identifies a (specific) type that is covered by or descended from the target type.
 - 43 • For each discriminant of the target type that corresponds to a discriminant of the operand type, its value is that of the corresponding discriminant of the operand value; if it

- corresponds to more than one discriminant of the operand type, a check is made that all these discriminants are equal in the operand value. 44
- For each discriminant of the target type that corresponds to a discriminant that is specified by the `derived_type_definition` for some ancestor of the operand type (or if class-wide, some ancestor of the specific type identified by the tag of the operand), its value in the result is that specified by the `derived_type_definition`. 45
 - For each discriminant of the operand type that corresponds to a discriminant that is specified by the `derived_type_definition` for some ancestor of the target type, a check is made that in the operand value it equals the value specified for it. 45
 - For each discriminant of the result, a check is made that its value belongs to its subtype. 46
- Access Type Conversion 47
- For an access-to-object type, a check is made that the accessibility level of the operand type is not deeper than that of the target type. 48
 - If the operand value is null, the result of the conversion is the null value of the target type. 49/2
 - If the operand value is not null, then the result designates the same object (or subprogram) as is designated by the operand value, but viewed as being of the target designated subtype (or profile); any checks associated with evaluating a conversion to the target designated subtype are performed. 50
- After conversion of the value to the target type, if the target subtype is constrained, a check is performed that the value satisfies this constraint. If the target subtype excludes null, then a check is made that the value is not null. 51/2
- For the evaluation of a view conversion, the operand `name` is evaluated, and a new view of the object denoted by the operand is created, whose type is the target type; if the target type is composite, checks are performed as above for a value conversion. 52
- The properties of this new view are as follows: 53
- If the target type is composite, the bounds or discriminants (if any) of the view are as defined above for a value conversion; each nondiscriminant component of the view denotes the matching component of the operand object; the subtype of the view is constrained if either the target subtype or the operand object is constrained, or if the target subtype is indefinite, or if the operand type is a descendant of the target type and has discriminants that were not inherited from the target type; 54/1
 - If the target type is tagged, then an assignment to the view assigns to the corresponding part of the object denoted by the operand; otherwise, an assignment to the view assigns to the object, after converting the assigned value to the subtype of the object (which might raise `Constraint_Error`); 55
 - Reading the value of the view yields the result of converting the value of the operand object to the target subtype (which might raise `Constraint_Error`), except if the object is of an access type and the view conversion is passed as an `out` parameter; in this latter case, the value of the operand object is used to initialize the formal parameter without checking against any constraint of the target subtype (see 6.4.1). 56
- If an `Accessibility_Check` fails, `Program_Error` is raised. Any other check associated with a conversion raises `Constraint_Error` if it fails. 57
- Conversion to a type is the same as conversion to an unconstrained subtype of the type. 58

NOTES

- 59 19 In addition to explicit `type_conversions`, type conversions are performed implicitly in situations where the expected type and the actual type of a construct differ, as is permitted by the type resolution rules (see 8.6). For example, an integer literal is of the type `universal_integer`, and is implicitly converted when assigned to a target of some specific integer type. Similarly, an actual parameter of a specific tagged type is implicitly converted when the corresponding formal parameter is of a class-wide type.
- 60 Even when the expected and actual types are the same, implicit subtype conversions are performed to adjust the array bounds (if any) of an operand to match the desired target subtype, or to raise `Constraint_Error` if the (possibly adjusted) value does not satisfy the constraints of the target subtype.
- 61/2 20 A ramification of the overload resolution rules is that the operand of an (explicit) `type_conversion` cannot be an allocator, an aggregate, a `string_literal`, a `character_literal`, or an `attribute_reference` for an `Access` or `Unchecked_Access` attribute. Similarly, such an expression enclosed by parentheses is not allowed. A `qualified_expression` (see 4.7) can be used instead of such a `type_conversion`.
- 62 21 The constraint of the target subtype has no effect for a `type_conversion` of an elementary type passed as an `out` parameter. Hence, it is recommended that the first subtype be specified as the target to minimize confusion (a similar recommendation applies to renaming and generic formal `in out` objects).

Examples

- 63 Examples of numeric type conversion:

```
64    Real(2**J)       -- value is converted to floating point
    Integer(1..6)     -- value is 2
    Integer(-0..4)    -- value is 0
```

- 65 Example of conversion between derived types:

```
66    type A_Form is new B_Form;
    X : A_Form;
    Y : B_Form;
67    X := A_Form(Y);
    Y := B_Form(X); -- the reverse conversion
```

- 69 Examples of conversions between array types:

```
70    type Sequence is array (Integer range <>) of Integer;
    subtype Dozen is Sequence(1 .. 12);
    Ledger : array(1 .. 100) of Integer;
71    Sequence(Ledger)   -- bounds are those of Ledger
    Sequence(Ledger(31 .. 42)) -- bounds are 31 and 42
    Dozen(Ledger(31 .. 42))   -- bounds are those of Dozen
```

4.7 Qualified Expressions

- 1 A `qualified_expression` is used to state explicitly the type, and to verify the subtype, of an operand that is either an expression or an aggregate.

Syntax

- 2 `qualified_expression ::=`
 `subtype_mark'(expression) | subtype_mark'aggregate`

Name Resolution Rules

- 3 The *operand* (the `expression` or `aggregate`) shall resolve to be of the type determined by the `subtype_mark`, or a universal type that covers it.

Dynamic Semantics

The evaluation of a `qualified_expression` evaluates the operand (and if of a universal type, converts it to the type determined by the `subtype_mark`) and checks that its value belongs to the subtype denoted by the `subtype_mark`. The exception `Constraint_Error` is raised if this check fails.

NOTES

22 When a given context does not uniquely identify an expected type, a `qualified_expression` can be used to do so. In particular, if an overloaded name or aggregate is passed to an overloaded subprogram, it might be necessary to qualify the operand to resolve its type.

Examples

Examples of disambiguating expressions using qualification:

```
type Mask is (Fix, Dec, Exp, Signif);
type Code is (Fix, Cla, Dec, Tnz, Sub);
Print (Mask' (Dec));    -- Dec is of type Mask
Print (Code' (Dec));   -- Dec is of type Code
for J in Code' (Fix) .. Code' (Dec) loop ... -- qualification needed for either Fix or Dec
for J in Code range Fix .. Dec loop ...      -- qualification unnecessary
for J in Code' (Fix) .. Dec loop ...         -- qualification unnecessary for Dec
Dozen' (1 | 3 | 5 | 7 => 2, others => 0) -- see 4.6
```

4.8 Allocators

The evaluation of an allocator creates an object and yields an access value that designates the object.

Syntax

```
allocator ::= new subtype_indication | new qualified_expression
```

Name Resolution Rules

The expected type for an allocator shall be a single access-to-object type with designated type *D* such that either *D* covers the type determined by the `subtype_mark` of the `subtype_indication` or `qualified_expression`, or the expected type is anonymous and the determined type is *D*'Class.

Legality Rules

An *initialized* allocator is an allocator with a `qualified_expression`. An *uninitialized* allocator is one with a `subtype_indication`. In the `subtype_indication` of an uninitialized allocator, a `constraint` is permitted only if the `subtype_mark` denotes an unconstrained composite subtype; if there is no `constraint`, then the `subtype_mark` shall denote a definite subtype.

If the type of the allocator is an access-to-constant type, the allocator shall be an initialized allocator.

If the designated type of the type of the allocator is class-wide, the accessibility level of the type determined by the `subtype_indication` or `qualified_expression` shall not be statically deeper than that of the type of the allocator.

If the designated subtype of the type of the allocator has one or more unconstrained access discriminants, then the accessibility level of the anonymous access type of each access discriminant, as determined by the `subtype_indication` or `qualified_expression` of the allocator, shall not be statically deeper than that of the type of the allocator (see 3.10.2).

- 5.3/2 An allocator shall not be of an access type for which the Storage_Size has been specified by a static expression with value zero or is defined by the language to be zero. In addition to the places where Legality Rules normally apply (see 12.3), this rule applies also in the private part of an instance of a generic unit. This rule does not apply in the body of a generic unit or within a body declared within the declarative region of a generic unit, if the type of the allocator is a descendant of a formal access type declared within the formal part of the generic unit.

Static Semantics

- 6/2 If the designated type of the type of the allocator is elementary, then the subtype of the created object is the designated subtype. If the designated type is composite, then the subtype of the created object is the designated subtype when the designated subtype is constrained or there is a partial view of the designated type that is constrained; otherwise, the created object is constrained by its initial value (even if the designated subtype is unconstrained with defaults).

Dynamic Semantics

- 7/2 For the evaluation of an initialized allocator, the evaluation of the `qualified_expression` is performed first. An object of the designated type is created and the value of the `qualified_expression` is converted to the designated subtype and assigned to the object.
- 8 For the evaluation of an uninitialized allocator, the elaboration of the `subtype_indication` is performed first. Then:
- 9/2 • If the designated type is elementary, an object of the designated subtype is created and any implicit initial value is assigned;
 - 10/2 • If the designated type is composite, an object of the designated type is created with tag, if any, determined by the `subtype_mark` of the `subtype_indication`. This object is then initialized by default (see 3.3.1) using the `subtype_indication` to determine its nominal subtype. A check is made that the value of the object belongs to the designated subtype. `Constraint_Error` is raised if this check fails. This check and the initialization of the object are performed in an arbitrary order.
- 10.1/2 For any allocator, if the designated type of the type of the allocator is class-wide, then a check is made that the accessibility level of the type determined by the `subtype_indication`, or by the tag of the value of the `qualified_expression`, is not deeper than that of the type of the allocator. If the designated subtype of the allocator has one or more unconstrained access discriminants, then a check is made that the accessibility level of the anonymous access type of each access discriminant is not deeper than that of the type of the allocator. `Program_Error` is raised if either such check fails.
- 10.2/2 If the object to be created by an allocator has a controlled or protected part, and the finalization of the collection of the type of the allocator (see 7.6.1) has started, `Program_Error` is raised.
- 10.3/2 If the object to be created by an allocator contains any tasks, and the master of the type of the allocator is completed, and all of the dependent tasks of the master are terminated (see 9.3), then `Program_Error` is raised.
- 11 If the created object contains any tasks, they are activated (see 9.2). Finally, an access value that designates the created object is returned.

Bounded (Run-Time) Errors

- 11.1/2 It is a bounded error if the finalization of the collection of the type (see 7.6.1) of the allocator has started. If the error is detected, `Program_Error` is raised. Otherwise, the allocation proceeds normally.

NOTES

- 23 Allocators cannot create objects of an abstract type. See 3.9.3. 12
- 24 If any part of the created object is controlled, the initialization includes calls on corresponding Initialize or Adjust procedures. See 7.6. 13
- 25 As explained in 13.11, “Storage Management”, the storage for an object allocated by an allocator comes from a storage pool (possibly user defined). The exception Storage_Error is raised by an allocator if there is not enough storage. Instances of Unchecked_Deallocation may be used to explicitly reclaim storage. 14
- 26 Implementations are permitted, but not required, to provide garbage collection (see 13.11.3). 15

*Examples**Examples of allocators:*

```

new Cell'(0, null, null)          -- initialized explicitly, see 3.10.1
new Cell'(Value => 0, Succ => null, Pred => null) -- initialized explicitly
new Cell                           -- not initialized

new Matrix(1 .. 10, 1 .. 20)       -- the bounds only are given
new Matrix'(1 .. 10 => (1 .. 20 => 0.0)) -- initialized explicitly

new Buffer(100)                   -- the discriminant only is given
new Buffer'(Size => 80, Pos => 0, Value => (1 .. 80 => 'A')) -- initialized explicitly

Expr_Ptr'(new Literal)           -- allocator for access-to-class-wide type, see 3.9.1
Expr_Ptr'(new Literal'(Expression with 3.5)) -- initialized explicitly

```

4.9 Static Expressions and Static Subtypes

Certain expressions of a scalar or string type are defined to be static. Similarly, certain discrete ranges are defined to be static, and certain scalar and string subtypes are defined to be static subtypes. *Static* means determinable at compile time, using the declared properties or values of the program entities. 1

A static expression is a scalar or string expression that is one of the following: 2

- a numeric_literal; 3
- a string_literal of a static string subtype; 4
- a name that denotes the declaration of a named number or a static constant; 5
- a function_call whose function_name or function_prefix statically denotes a static function, and whose actual parameters, if any (whether given explicitly or by default), are all static expressions; 6
- an attribute_reference that denotes a scalar value, and whose prefix denotes a static scalar subtype; 7
- an attribute_reference whose prefix statically denotes a statically constrained array object or array subtype, and whose attribute_designator is First, Last, or Length, with an optional dimension; 8
- a type_conversion whose subtype_mark denotes a static scalar subtype, and whose operand is a static expression; 9
- a qualified_expression whose subtype_mark denotes a static (scalar or string) subtype, and whose operand is a static expression; 10
- a membership test whose simple_expression is a static expression, and whose range is a static range or whose subtype_mark denotes a static (scalar or string) subtype; 11
- a short-circuit control form both of whose relations are static expressions; 12
- a static expression enclosed in parentheses. 13

- 14 A name *statically denotes* an entity if it denotes the entity and:
- 15 • It is a `direct_name`, expanded name, or `character_literal`, and it denotes a declaration other than a `renaming_declaration`; or
 - 16 • It is an `attribute_reference` whose `prefix` statically denotes some entity; or
 - 17 • It denotes a `renaming_declaration` with a name that statically denotes the renamed entity.
- 18 A *static function* is one of the following:
- 19 • a predefined operator whose parameter and result types are all scalar types none of which are descendants of formal scalar types;
 - 20 • a predefined concatenation operator whose result type is a string type;
 - 21 • an enumeration literal;
 - 22 • a language-defined attribute that is a function, if the `prefix` denotes a static scalar subtype, and if the parameter and result types are scalar.
- 23 In any case, a generic formal subprogram is not a static function.
- 24 A *static constant* is a constant view declared by a full constant declaration or an `object_renaming_declaration` with a static nominal subtype, having a value defined by a static scalar expression or by a static string expression whose value has a length not exceeding the maximum length of a `string_literal` in the implementation.
- 25 A *static range* is a `range` whose bounds are static expressions, or a `range_attribute_reference` that is equivalent to such a `range`. A *static discrete_range* is one that is a static range or is a `subtype_indication` that defines a static scalar subtype. The base range of a scalar type is a static range, unless the type is a descendant of a formal scalar type.
- 26/2 A *static subtype* is either a *static scalar subtype* or a *static string subtype*. A static scalar subtype is an unconstrained scalar subtype whose type is not a descendant of a formal type, or a constrained scalar subtype formed by imposing a compatible static constraint on a static scalar subtype. A static string subtype is an unconstrained string subtype whose index subtype and component subtype are static, or a constrained string subtype formed by imposing a compatible static constraint on a static string subtype. In any case, the subtype of a generic formal object of mode `in out`, and the result subtype of a generic formal function, are not static.
- 27 The different kinds of *static constraint* are defined as follows:
- 28 • A null constraint is always static;
 - 29 • A scalar constraint is static if it has no `range_constraint`, or one with a static range;
 - 30 • An index constraint is static if each `discrete_range` is static, and each index subtype of the corresponding array type is static;
 - 31 • A discriminant constraint is static if each `expression` of the constraint is static, and the subtype of each discriminant is static.
- 31.1/2 In any case, the constraint of the first subtype of a scalar formal type is neither static nor null.
- 32 A subtype is *statically constrained* if it is constrained, and its constraint is static. An object is *statically constrained* if its nominal subtype is statically constrained, or if it is a static string constant.

Legality Rules

A static expression is evaluated at compile time except when it is part of the right operand of a static short-circuit control form whose value is determined by its left operand. This evaluation is performed exactly, without performing Overflow_Checks. For a static expression that is evaluated:

- The expression is illegal if its evaluation fails a language-defined check other than Overflow_Check. 33
- If the expression is not part of a larger static expression and the expression is expected to be of a single specific type, then its value shall be within the base range of its expected type. Otherwise, the value may be arbitrarily large or small. 35/2
- If the expression is of type *universal_real* and its expected type is a decimal fixed point type, then its value shall be a multiple of the *small* of the decimal type. This restriction does not apply if the expected type is a descendant of a formal scalar type (or a corresponding actual type in an instance). 36/2

In addition to the places where Legality Rules normally apply (see 12.3), the above restrictions also apply in the private part of an instance of a generic unit. 37/2

Implementation Requirements

For a real static expression that is not part of a larger static expression, and whose expected type is not a descendant of a formal type, the implementation shall round or truncate the value (according to the Machine_Rounds attribute of the expected type) to the nearest machine number of the expected type; if the value is exactly half-way between two machine numbers, the rounding performed is implementation-defined. If the expected type is a descendant of a formal type, or if the static expression appears in the body of an instance of a generic unit and the corresponding expression is nonstatic in the corresponding generic body, then no special rounding or truncating is required — normal accuracy rules apply (see Annex G). 38/2

Implementation Advice

For a real static expression that is not part of a larger static expression, and whose expected type is not a descendant of a formal type, the rounding should be the same as the default rounding for the target system. 38.1/2

NOTES

- 27 An expression can be static even if it occurs in a context where staticness is not required. 39
- 28 A static (or run-time) type_conversion from a real type to an integer type performs rounding. If the operand value is exactly half-way between two integers, the rounding is performed away from zero. 40

Examples

Examples of static expressions:

```
1 + 1      -- 2
abs(-10)*3 -- 30

Kilo : constant := 1000;
Mega : constant := Kilo*Kilo;    -- 1_000_000
Long : constant := Float'Digits*2;

Half_Pi   : constant := Pi/2;           -- see 3.3.2
Deg_To_Rad : constant := Half_Pi/90;
Rad_To_Deg : constant := 1.0/Deg_To_Rad; -- equivalent to 1.0/((3.14159_26536/2)/90)
```

4.9.1 Statically Matching Constraints and Subtypes

Static Semantics

- 1/2 A constraint *statically matches* another constraint if:
 - 1.1/2 • both are null constraints;
 - 1.2/2 • both are static and have equal corresponding bounds or discriminant values;
 - 1.3/2 • both are nonstatic and result from the same elaboration of a constraint of a `subtype_indication` or the same evaluation of a `range` of a `discrete_subtype_definition`; or
 - 1.4/2 • both are nonstatic and come from the same `formal_type_declaration`.
- 2/2 A subtype *statically matches* another subtype of the same type if they have statically matching constraints, and, for access subtypes, either both or neither exclude null. Two anonymous access-to-object subtypes statically match if their designated subtypes statically match, and either both or neither exclude null, and either both or neither are access-to-constant. Two anonymous access-to-subprogram subtypes statically match if their designated profiles are subtype conformant, and either both or neither exclude null.
- 3 Two ranges of the same type *statically match* if both result from the same evaluation of a `range`, or if both are static and have equal corresponding bounds.
- 4 A constraint is *statically compatible* with a scalar subtype if it statically matches the constraint of the subtype, or if both are static and the constraint is compatible with the subtype. A constraint is *statically compatible* with an access or composite subtype if it statically matches the constraint of the subtype, or if the subtype is unconstrained. One subtype is *statically compatible* with a second subtype if the constraint of the first is statically compatible with the second subtype.

Section 5: Statements

A statement defines an action to be performed upon its execution.

This section describes the general rules applicable to all statements. Some statements are discussed in later sections: `Procedure_call_statements` and `return_statements` are described in 6, “Subprograms”. `Entry_call_statements`, `requeue_statements`, `delay_statements`, `accept_statements`, `select_statements`, and `abort_statements` are described in 9, “Tasks and Synchronization”. `Raise_statements` are described in 11, “Exceptions”, and `code_statements` in 13. The remaining forms of statements are presented in this section.

5.1 Simple and Compound Statements - Sequences of Statements

A statement is either simple or compound. A **simple_statement** encloses no other statement. A **compound_statement** can enclose simple statements and other compound statements.

Syntax

```

sequence_of_statements ::= statement {statement}

statement ::=

  {label} simple_statement | {label} compound_statement

simple_statement ::= null_statement

  | assignment_statement           | exit_statement
  | goto_statement                | procedure_call_statement
  | simple_return_statement       | entry_call_statement
  | requeue_statement            | delay_statement
  | abort_statement              | raise_statement
  | code_statement

compound_statement ::=

  if_statement                   | case_statement
  | loop_statement               | block_statement
  | extended_return_statement    |
  | accept_statement             | select_statement

null_statement ::= null;

label ::= <<label_statement_identifier>>

statement_identifier ::= direct_name

The direct_name of a statement identifier shall be an identifier (n)

```

Name Resolution Rules

The `direct_name` of a `statement_identifier` shall resolve to denote its corresponding implicit declaration (see below).

Legality Rules

Distinct identifiers shall be used for all **statement_identifiers** that appear in the same body, including inner block **statements** but excluding inner program units.

Static Semantics

- 12 For each **statement_identifier**, there is an implicit declaration (with the specified identifier) at the end of the **declarative_part** of the innermost **block_statement** or body that encloses the **statement_identifier**. The implicit declarations occur in the same order as the **statement_identifiers** occur in the source text. If a usage name denotes such an implicit declaration, the entity it denotes is the **label**, **loop_statement**, or **block_statement** with the given **statement_identifier**.

Dynamic Semantics

- 13 The execution of a **null_statement** has no effect.
- 14/2 A *transfer of control* is the run-time action of an **exit_statement**, return statement, **goto_statement**, or **queue_statement**, selection of a **terminate_alternative**, raising of an exception, or an abort, which causes the next action performed to be one other than what would normally be expected from the other rules of the language. As explained in 7.6.1, a transfer of control can cause the execution of constructs to be completed and then left, which may trigger finalization.
- 15 The execution of a **sequence_of_statements** consists of the execution of the individual **statements** in succession until the **sequence_is** completed.

NOTES

- 16 1 A **statement_identifier** that appears immediately within the declarative region of a named **loop_statement** or an **accept_statement** is nevertheless implicitly declared immediately within the declarative region of the innermost enclosing body or **block_statement**; in other words, the expanded name for a named statement is not affected by whether the statement occurs inside or outside a named loop or an **accept_statement** — only nesting within **block_statements** is relevant to the form of its expanded name.

Examples

- 17 Examples of labeled statements:

```
<<Here>> <<Ici>> <<Aquici>> <<Hier>> null;
<<After>> X := 1;
```

5.2 Assignment Statements

- 1 An **assignment_statement** replaces the current value of a variable with the result of evaluating an expression.

Syntax

- 2 **assignment_statement ::=**
 variable_name := expression;
- 3 The execution of an **assignment_statement** includes the evaluation of the **expression** and the *assignment* of the value of the **expression** into the *target*. An assignment operation (as opposed to an **assignment_statement**) is performed in other contexts as well, including object initialization and by-copy parameter passing. The *target* of an assignment operation is the view of the object to which a value is being assigned; the target of an **assignment_statement** is the variable denoted by the **variable_name**.

Name Resolution Rules

- 4/2 The **variable_name** of an **assignment_statement** is expected to be of any type. The expected type for the **expression** is the type of the target.

Legality Rules

The target denoted by the *variable_name* shall be a variable of a nonlimited type.

5/2

If the target is of a tagged class-wide type *TClass*, then the **expression** shall either be dynamically tagged, or of type *T* and tag-indeterminate (see 3.9.2).

6

Dynamic Semantics

For the execution of an **assignment_statement**, the *variable_name* and the **expression** are first evaluated in an arbitrary order.

7

When the type of the target is class-wide:

8

- If the **expression** is tag-indeterminate (see 3.9.2), then the controlling tag value for the **expression** is the tag of the target;
- Otherwise (the **expression** is dynamically tagged), a check is made that the tag of the value of the **expression** is the same as that of the target; if this check fails, *Constraint_Error* is raised.

9

10

The value of the **expression** is converted to the subtype of the target. The conversion might raise an exception (see 4.6).

11

In cases involving controlled types, the target is finalized, and an anonymous object might be used as an intermediate in the assignment, as described in 7.6.1, “Completion and Finalization”. In any case, the converted value of the **expression** is then *assigned* to the target, which consists of the following two steps:

12

- The value of the target becomes the converted value.
- If any part of the target is controlled, its value is adjusted as explained in clause 7.6.

13

14

NOTES

2 The tag of an object never changes; in particular, an **assignment_statement** does not change the tag of the target.

15

This paragraph was deleted.

16/2

Examples

Examples of assignment statements:

17

```
Value := Max_Value - 1;
Shade := Blue;
Next_Frame(F)(M, N) := 2.5;           -- see 4.1.1
U := Dot_Product(V, W);              -- see 6.3
Writer := (Status => Open, Unit => Printer, Line_Count => 60); -- see 3.8.1
Next_Car.all := (72074, null);        -- see 3.10.1
```

18

19

20

Examples involving scalar subtype conversions:

21

```
I, J : Integer range 1 .. 10 := 5;
K     : Integer range 1 .. 20 := 15;
...
I := J;   -- identical ranges
K := J;   -- compatible ranges
J := K;   -- will raise Constraint_Error if K > 10
```

22

23

Examples involving array subtype conversions:

24

```
A : String(1 .. 31);
B : String(3 .. 33);
...
A := B; -- same number of components
```

25

26

27 A(1 .. 9) := "tar sauce";
 A(4 .. 12) := A(1 .. 9); -- A(1 .. 12) = "tartar sauce"

NOTES

3 *Notes on the examples:* Assignment_statements are allowed even in the case of overlapping slices of the same array, because the variable_name and expression are both evaluated before copying the value into the variable. In the above example, an implementation yielding A[1 .. 12] = "tartartartar" would be incorrect.

5.3 If Statements

- 1 An `if_statement` selects for execution at most one of the enclosed `sequences_of_statements`, depending on the (truth) value of one or more corresponding conditions.

Syntax

```

if _statement ::=

  if condition then
    sequence_of_statements
  {elsif condition then
    sequence_of_statements}
  [else
    sequence_of_statements]
  end if;

condition ::= boolean_expression

```

Name Resolution Rules

- 4 A condition is expected to be of any boolean type.

Dynamic Semantics

- 5 For the execution of an `if_statement`, the condition specified after `if`, and any conditions specified after `elsif`, are evaluated in succession (treating a final `else` as `elsif True then`), until one evaluates to True or all conditions are evaluated and yield False. If a condition evaluates to True, then the corresponding sequence of statements is executed; otherwise none of them is executed.

Examples

- ## 6 Examples of if statements:

```

if Month = December and Day = 31 then
  Month := January;
  Day   := 1;
  Year  := Year + 1;
end if;

if Line_Too_Short then
  raise Layout_Error;
elsif Line_Full then
  New_Line;
  Put(Item);
else
  Put(Item);
end if;

if My_Car.Owner.Vehicle /= My_Car then
  Report ("Incorrect data");
end if;

```

-- see 3.10.1

5.4 Case Statements

A `case_statement` selects for execution one of a number of alternative `sequences_of_statements`; the chosen alternative is defined by the value of an expression.

1

Syntax

```
case_statement ::=  
  case expression is  
    case_statement_alternative  
    {case_statement_alternative}  
  end case;  
  
case_statement_alternative ::=  
  when discrete_choice_list =>  
    sequence_of_statements
```

2

3

Name Resolution Rules

The `expression` is expected to be of any discrete type. The expected type for each `discrete_choice` is the type of the `expression`.

Legality Rules

The expressions and `discrete_ranges` given as `discrete_choices` of a `case_statement` shall be static. A `discrete_choice others`, if present, shall appear alone and in the last `discrete_choice_list`.

The possible values of the `expression` shall be covered as follows:

- If the `expression` is a name (including a `type_conversion` or a `function_call`) having a static and constrained nominal subtype, or is a `qualified_expression` whose `subtype_mark` denotes a static and constrained scalar subtype, then each non-`others` `discrete_choice` shall cover only values in that subtype, and each value of that subtype shall be covered by some `discrete_choice` (either explicitly or by `others`).
- If the type of the `expression` is `root_integer`, `universal_integer`, or a descendant of a formal scalar type, then the `case_statement` shall have an `others` `discrete_choice`.
- Otherwise, each value of the base range of the type of the `expression` shall be covered (either explicitly or by `others`).

Two distinct `discrete_choices` of a `case_statement` shall not cover the same value.

Dynamic Semantics

For the execution of a `case_statement` the `expression` is first evaluated.

If the value of the `expression` is covered by the `discrete_choice_list` of some `case_statement_alternative`, then the `sequence_of_statements` of the `_alternative` is executed.

Otherwise (the value is not covered by any `discrete_choice_list`, perhaps due to being outside the base range), `Constraint_Error` is raised.

NOTES

- 4 The execution of a `case_statement` chooses one and only one alternative. Qualification of the `expression` of a `case_statement` by a static subtype can often be used to limit the number of choices that need be given explicitly.

Examples

15 Examples of case statements:

```

16    case Sensor is
        when Elevation => Record_Elevation(Sensor_Value);
        when Azimuth    => Record_Azimuth  (Sensor_Value);
        when Distance   => Record_Distance (Sensor_Value);
        when others      => null;
    end case;

17    case Today is
        when Mon         => Compute_Initial_Balance;
        when Fri         => Compute_Closing_Balance;
        when Tue .. Thu => Generate_Report(Today);
        when Sat .. Sun => null;
    end case;

18    case Bin_Number(Count) is
        when 1          => Update_Bin(1);
        when 2          => Update_Bin(2);
        when 3 | 4 =>
            Empty_Bin(1);
            Empty_Bin(2);
        when others     => raise Error;
    end case;

```

5.5 Loop Statements

- 1 A loop_statement includes a sequence_of_statements that is to be executed repeatedly, zero or more times.

Syntax

```

2      loop_statement ::= 
3          [loop_statement_identifier:] 
4              [iteration_scheme] loop
5                  sequence_of_statements
6                  end loop [loop_identifier];
7
8      iteration_scheme ::= while condition
9          | for loop_parameter_specification
10
11     loop_parameter_specification ::= 
12         defining_identifier in [reverse] discrete_subtype_definition
13
14     If a loop_statement has a loop_statement_identifier, then the identifier shall be repeated after the end loop; otherwise, there shall not be an identifier after the end loop.

```

Static Semantics

- 6 A loop_parameter_specification declares a *loop parameter*, which is an object whose subtype is that defined by the discrete_subtype_definition.

Dynamic Semantics

- 7 For the execution of a loop_statement, the sequence_of_statements is executed repeatedly, zero or more times, until the loop_statement is complete. The loop_statement is complete when a transfer of control occurs that transfers control out of the loop, or, in the case of an iteration_scheme, as specified below.
- 8 For the execution of a loop_statement with a **while** iteration_scheme, the condition is evaluated before each execution of the sequence_of_statements; if the value of the condition is True, the sequence_of_statements is executed; if False, the execution of the loop_statement is complete.

For the execution of a `loop_statement` with a `for` iteration_scheme, the `loop_parameter_specification` is first elaborated. This elaboration creates the loop parameter and elaborates the `discrete_subtype_definition`. If the `discrete_subtype_definition` defines a subtype with a null range, the execution of the `loop_statement` is complete. Otherwise, the `sequence_of_statements` is executed once for each value of the discrete subtype defined by the `discrete_subtype_definition` (or until the loop is left as a consequence of a transfer of control). Prior to each such iteration, the corresponding value of the discrete subtype is assigned to the loop parameter. These values are assigned in increasing order unless the reserved word `reverse` is present, in which case the values are assigned in decreasing order.

NOTES

- 5 A loop parameter is a constant; it cannot be updated within the `sequence_of_statements` of the loop (see 3.3).
- 6 An object_declaration should not be given for a loop parameter, since the loop parameter is automatically declared by the `loop_parameter_specification`. The scope of a loop parameter extends from the `loop_parameter_specification` to the end of the `loop_statement`, and the visibility rules are such that a loop parameter is only visible within the `sequence_of_statements` of the loop.
- 7 The `discrete_subtype_definition` of a `for` loop is elaborated just once. Use of the reserved word `reverse` does not alter the discrete subtype defined, so that the following iteration_schemes are not equivalent; the first has a null range.

```
for J in reverse 1 .. 0
for J in 0 .. 1
```

Examples

Example of a loop statement without an iteration scheme:

```
loop
    Get(Current_Character);
    exit when Current_Character = '*';
end loop;
```

Example of a loop statement with a `while` iteration scheme:

```
while Bid(N).Price < Cut_Off.Price loop
    Record_Bid(Bid(N).Price);
    N := N + 1;
end loop;
```

Example of a loop statement with a `for` iteration scheme:

```
for J in Buffer'Range loop      -- works even with a null range
    if Buffer(J) /= Space then
        Put(Buffer(J));
    end if;
end loop;
```

Example of a loop statement with a name:

```
Summation:
  while Next /= Head loop      -- see 3.10.1
    Sum := Sum + Next.Value;
    Next := Next.Succ;
  end loop Summation;
```

5.6 Block Statements

- 1 A `block_statement` encloses a `handled_sequence_of_statements` optionally preceded by a `declarative_part`.

Syntax

```
2   block_statement ::=  
      [block_statement_identifier:]  
      [declare  
       declarative_part]  
      begin  
       handled_sequence_of_statements  
      end [block_identifier];
```

- 3 If a `block_statement` has a `block_statement_identifier`, then the identifier shall be repeated after the `end`; otherwise, there shall not be an identifier after the `end`.

Static Semantics

- 4 A `block_statement` that has no explicit `declarative_part` has an implicit empty `declarative_part`.

Dynamic Semantics

- 5 The execution of a `block_statement` consists of the elaboration of its `declarative_part` followed by the execution of its `handled_sequence_of_statements`.

Examples

- 6 Example of a `block statement` with a local variable:

```
7   Swap:  
     declare  
       Temp : Integer;  
     begin  
       Temp := V; V := U; U := Temp;  
     end Swap;
```

5.7 Exit Statements

- 1 An `exit_statement` is used to complete the execution of an enclosing `loop_statement`; the completion is conditional if the `exit_statement` includes a condition.

Syntax

```
2   exit_statement ::=  
     exit [loop_name] [when condition];
```

Name Resolution Rules

- 3 The `loop_name`, if any, in an `exit_statement` shall resolve to denote a `loop_statement`.

Legality Rules

- 4 Each `exit_statement` applies to a `loop_statement`; this is the `loop_statement` being exited. An `exit_statement` with a name is only allowed within the `loop_statement` denoted by the name, and applies to that `loop_statement`. An `exit_statement` without a name is only allowed within a `loop_statement`, and applies to the innermost enclosing one. An `exit_statement` that applies to a given `loop_statement` shall not

appear within a body or `accept_statement`, if this construct is itself enclosed by the given `loop_statement`.

Dynamic Semantics

For the execution of an `exit_statement`, the condition, if present, is first evaluated. If the value of the condition is True, or if there is no condition, a transfer of control is done to complete the `loop_statement`. If the value of the condition is False, no transfer of control takes place.

NOTES

- 8 Several nested loops can be exited by an `exit_statement` that names the outer loop.

5

6

Examples

Examples of loops with exit statements:

```
for N in 1 .. Max_Num_Items loop
    Get_New_Item(New_Item);
    Merge_Item(New_Item, Storage_File);
    exit when New_Item = Terminal_Item;
end loop;

Main_Cycle:
loop
    -- initial statements
    exit Main_Cycle when Found;
    -- final statements
end loop Main_Cycle;
```

7

8

9

5.8 Goto Statements

A `goto_statement` specifies an explicit transfer of control from this `statement` to a target statement with a given label.

1

Syntax

`goto_statement ::= goto label_name;`

2

Name Resolution Rules

The `label_name` shall resolve to denote a `label`; the `statement` with that `label` is the *target statement*.

3

Legality Rules

The innermost `sequence_of_statements` that encloses the target statement shall also enclose the `goto_statement`. Furthermore, if a `goto_statement` is enclosed by an `accept_statement` or a body, then the target statement shall not be outside this enclosing construct.

4

Dynamic Semantics

The execution of a `goto_statement` transfers control to the target statement, completing the execution of any `compound_statement` that encloses the `goto_statement` but does not enclose the target.

5

NOTES

- 9 The above rules allow transfer of control to a `statement` of an enclosing `sequence_of_statements` but not the reverse. Similarly, they prohibit transfers of control such as between alternatives of a `case_statement`, `if_statement`, or `select_statement`; between `exception_handlers`; or from an `exception_handler` of a `handled_sequence_of_statements` back to its `sequence_of_statements`.

6

Examples

7 Example of a loop containing a goto statement:

```
8    <>Sort>>
  for I in 1 .. N-1 loop
    if A(I) > A(I+1) then
      Exchange(A(I), A(I+1));
      goto Sort;
    end if;
  end loop;
```

Section 6: Subprograms

A subprogram is a program unit or intrinsic operation whose execution is invoked by a subprogram call. There are two forms of subprogram: procedures and functions. A procedure call is a statement; a function call is an expression and returns a value. The definition of a subprogram can be given in two parts: a subprogram declaration defining its interface, and a subprogram_body defining its execution. Operators and enumeration literals are functions.

A *callable entity* is a subprogram or entry (see Section 9). A callable entity is invoked by a *call*; that is, a subprogram call or entry call. A *callable construct* is a construct that defines the action of a call upon a callable entity: a subprogram_body, entry_body, or accept_statement.

6.1 Subprogram Declarations

A subprogram_declaration declares a procedure or function.

<i>Syntax</i>	
subprogram_declaration ::=	2/2
[overriding_indicator]	
subprogram_specification;	
<i>This paragraph was deleted.</i>	3/2
subprogram_specification ::=	4/2
procedure_specification	
function_specification	
procedure_specification ::= procedure defining_program_unit_name parameter_profile	4.1/2
function_specification ::= function defining_designator parameter_and_result_profile	4.2/2
designator ::= [parent_unit_name .]identifier operator_symbol	5
defining_designator ::= defining_program_unit_name defining_operator_symbol	6
defining_program_unit_name ::= [parent_unit_name .]defining_identifier	7
The optional parent_unit_name is only allowed for library units (see 10.1.1).	8
operator_symbol ::= string_literal	9
The sequence of characters in an operator_symbol shall form a reserved word, a delimiter, or compound delimiter that corresponds to an operator belonging to one of the six categories of operators defined in clause 4.5.	10/2
defining_operator_symbol ::= operator_symbol	11
parameter_profile ::= [formal_part]	12
parameter_and_result_profile ::=	13/2
[formal_part] return [null_exclusion] subtype_mark	
[formal_part] access_definition	
formal_part ::=	14
(parameter_specification {; parameter_specification})	
parameter_specification ::=	15/2
defining_identifier_list : mode [null_exclusion] subtype_mark [: default_expression]	
defining_identifier_list : access_definition [: default_expression]	
mode ::= [in] in out out	16

Name Resolution Rules

- 17 A *formal parameter* is an object directly visible within a `subprogram_body` that represents the actual parameter passed to the subprogram in a call; it is declared by a `parameter_specification`. For a formal parameter, the expected type for its `default_expression`, if any, is that of the formal parameter.

Legality Rules

- 18 The *parameter mode* of a formal parameter conveys the direction of information transfer with the actual parameter: `in`, `in out`, or `out`. Mode `in` is the default, and is the mode of a parameter defined by an `access_definition`. The formal parameters of a function, if any, shall have the mode `in`.
- 19 A `default_expression` is only allowed in a `parameter_specification` for a formal parameter of mode `in`.
- 20/2 A `subprogram_declaration` or a `generic_subprogram_declaration` requires a completion: a body, a `renaming_declaration` (see 8.5), or a `pragma Import` (see B.1). A completion is not allowed for an `abstract_subprogram_declaration` (see 3.9.3) or a `null_procedure_declaration` (see 6.7).
- 21 A name that denotes a formal parameter is not allowed within the `formal_part` in which it is declared, nor within the `formal_part` of a corresponding body or `accept_statement`.

Static Semantics

- 22 The *profile* of (a view of) a callable entity is either a `parameter_profile` or `parameter_and_result_profile`; it embodies information about the interface to that entity — for example, the profile includes information about parameters passed to the callable entity. All callable entities have a profile — enumeration literals, other subprograms, and entries. An access-to-subprogram type has a designated profile. Associated with a profile is a calling convention. A `subprogram_declaration` declares a procedure or a function, as indicated by the initial reserved word, with name and profile as given by its specification.
- 23/2 The nominal subtype of a formal parameter is the subtype determined by the optional `null_exclusion` and the `subtype_mark`, or defined by the `access_definition`, in the `parameter_specification`. The nominal subtype of a function result is the subtype determined by the optional `null_exclusion` and the `subtype_mark`, or defined by the `access_definition`, in the `parameter_and_result_profile`.
- 24/2 An *access parameter* is a formal `in` parameter specified by an `access_definition`. An *access result type* is a function result type specified by an `access_definition`. An access parameter or result type is of an anonymous access type (see 3.10). Access parameters of an access-to-object type allow dispatching calls to be controlled by access values. Access parameters of an access-to-subprogram type permit calls to subprograms passed as parameters irrespective of their accessibility level.
- 25 The *subtypes of a profile* are:
- 26 • For any non-access parameters, the nominal subtype of the parameter.
- 27/2
- For any access parameters of an access-to-object type, the designated subtype of the parameter type.
- 27.1/2
- For any access parameters of an access-to-subprogram type, the subtypes of the profile of the parameter type.
- 28/2
- For any non-access result, the nominal subtype of the function result.
- 28.1/2
- For any access result type of an access-to-object type, the designated subtype of the result type.
- 28.2/2
- For any access result type of an access-to-subprogram type, the subtypes of the profile of the result type.
- 29 The *types of a profile* are the types of those subtypes.

A subprogram declared by an `abstract_subprogram_declaration` is abstract; a subprogram declared by a `subprogram_declaration` is not. See 3.9.3, “Abstract Types and Subprograms”. Similarly, a procedure defined by a `null_procedure_declaration` is a null procedure; a procedure declared by a `subprogram_declaration` is not. See 6.7, “Null Procedures”.

An `overriding_indicator` is used to indicate whether overriding is intended. See 8.3.1, “Overriding Indicators”.

Dynamic Semantics

The elaboration of a `subprogram_declaration` has no effect.

NOTES

- 1 A `parameter_specification` with several identifiers is equivalent to a sequence of single `parameter_specifications`, as explained in 3.3.
- 2 Abstract subprograms do not have bodies, and cannot be used in a nondispatching call (see 3.9.3, “Abstract Types and Subprograms”).
- 3 The evaluation of `default_expressions` is caused by certain calls, as described in 6.4.1. They are not evaluated during the elaboration of the subprogram declaration.
- 4 Subprograms can be called recursively and can be called concurrently from multiple tasks.

Examples

Examples of subprogram declarations:

```
procedure Traverse_Tree;
procedure Increment(X : in out Integer);
procedure Right_Indent(Margin : out Line_Size);           -- see 3.5.4
procedure Switch(From, To : in out Link);                  -- see 3.10.1
function Random return Probability;                      -- see 3.5.7
function Min_Cell(X : Link) return Cell;                 -- see 3.10.1
function Next_Frame(K : Positive) return Frame;          -- see 3.10
function Dot_Product(Left, Right : Vector) return Real;   -- see 3.6
function "*" (Left, Right : Matrix) return Matrix;        -- see 3.6
```

Examples of `in` parameters with default expressions:

```
procedure Print_Header(Pages : in Natural;
                      Header : in Line    := (1 .. Line'Last => ' ');
                      Center : in Boolean := True); -- see 3.6
```

6.2 Formal Parameter Modes

A `parameter_specification` declares a formal parameter of mode `in`, `in out`, or `out`.

Static Semantics

A parameter is passed either *by copy* or *by reference*. When a parameter is passed by copy, the formal parameter denotes a separate object from the actual parameter, and any information transfer between the two occurs only before and after executing the subprogram. When a parameter is passed by reference, the formal parameter denotes (a view of) the object denoted by the actual parameter; reads and updates of the formal parameter directly reference the actual parameter object.

A type is a *by-copy type* if it is an elementary type, or if it is a descendant of a private type whose full type is a by-copy type. A parameter of a by-copy type is passed by copy.

- 4 A type is a *by-reference type* if it is a descendant of one of the following:
- 5 • a tagged type;
 - 6 • a task or protected type;
 - 7 • a nonprivate type with the reserved word **limited** in its declaration;
 - 8 • a composite type with a subcomponent of a by-reference type;
 - 9 • a private type whose full type is a by-reference type.
- 10 A parameter of a by-reference type is passed by reference. Each value of a by-reference type has an associated object. For a parenthesized expression, **qualified_expression**, or **type_conversion**, this object is the one associated with the operand.
- 11 For parameters of other types, it is unspecified whether the parameter is passed by copy or by reference.

Bounded (Run-Time) Errors

- 12 If one **name** denotes a part of a formal parameter, and a second **name** denotes a part of a distinct formal parameter or an object that is not part of a formal parameter, then the two **names** are considered *distinct access paths*. If an object is of a type for which the parameter passing mechanism is not specified, then it is a bounded error to assign to the object via one access path, and then read the value of the object via a distinct access path, unless the first access path denotes a part of a formal parameter that no longer exists at the point of the second access (due to leaving the corresponding callable construct). The possible consequences are that **Program_Error** is raised, or the newly assigned value is read, or some old value of the object is read.

NOTES

- 13 5 A formal parameter of mode **in** is a constant view (see 3.3); it cannot be updated within the **subprogram_body**.

6.3 Subprogram Bodies

- 1 A **subprogram_body** specifies the execution of a subprogram.

Syntax

2/2 **subprogram_body ::=**
 [**overriding_indicator**]
 subprogram_specification **is**
 declarative_part
 begin
 handled_sequence_of_statements
 end [**designator**];

- 3 If a **designator** appears at the end of a **subprogram_body**, it shall repeat the **defining_designator** of the **subprogram_specification**.

Legality Rules

- 4 In contrast to other bodies, a **subprogram_body** need not be the completion of a previous declaration, in which case the body declares the subprogram. If the body is a completion, it shall be the completion of a **subprogram_declaration** or **generic_subprogram_declaration**. The profile of a **subprogram_body** that completes a declaration shall conform fully to that of the declaration.

Static Semantics

A `subprogram_body` is considered a declaration. It can either complete a previous declaration, or itself be the initial declaration of the subprogram.

Dynamic Semantics

The elaboration of a non-generic `subprogram_body` has no other effect than to establish that the subprogram can from then on be called without failing the `Elaboration_Check`.

The execution of a `subprogram_body` is invoked by a subprogram call. For this execution the `declarative_part` is elaborated, and the `handled_sequence_of_statements` is then executed.

Examples

Example of procedure body:

```
procedure Push(E : in Element_Type; S : in out Stack) is
begin
  if S.Index = S.Size then
    raise Stack_Overflow;
  else
    S.Index := S.Index + 1;
    S.Space(S.Index) := E;
  end if;
end Push;
```

Example of a function body:

```
function Dot_Product(Left, Right : Vector) return Real is
  Sum : Real := 0.0;
begin
  Check(Left'First = Right'First and Left'Last = Right'Last);
  for J in Left'Range loop
    Sum := Sum + Left(J)*Right(J);
  end loop;
  return Sum;
end Dot_Product;
```

6.3.1 Conformance Rules

When subprogram profiles are given in more than one place, they are required to conform in one of four ways: type conformance, mode conformance, subtype conformance, or full conformance.

Static Semantics

As explained in B.1, “Interfacing Pragmas”, a *convention* can be specified for an entity. Unless this International Standard states otherwise, the default convention of an entity is Ada. For a callable entity or access-to-subprogram type, the convention is called the *calling convention*. The following conventions are defined by the language:

- The default calling convention for any subprogram not listed below is *Ada*. A `pragma Convention`, `Import`, or `Export` may be used to override the default calling convention (see B.1).
- The *Intrinsic* calling convention represents subprograms that are “built in” to the compiler. The default calling convention is Intrinsic for the following:
 - an enumeration literal;
 - a “/=” operator declared implicitly due to the declaration of “=” (see 6.6);
 - any other implicitly declared subprogram unless it is a dispatching operation of a tagged type;

- 8 • an inherited subprogram of a generic formal tagged type with unknown discriminants;
 - 9 • an attribute that is a subprogram;
 - 10/2 • a subprogram declared immediately within a **protected_body**;
 - 10.1/2 • any prefixed view of a subprogram (see 4.1.3).
- 11 The Access attribute is not allowed for Intrinsic subprograms.
- 12 • The default calling convention is *protected* for a protected subprogram, and for an access-to-subprogram type with the reserved word **protected** in its definition.
 - 13 • The default calling convention is *entry* for an entry.
- 13.1/2 • The calling convention for an anonymous access-to-subprogram parameter or anonymous access-to-subprogram result is *protected* if the reserved word **protected** appears in its definition and otherwise is the convention of the subprogram that contains the parameter.
- 13.2/1 • If not specified above as Intrinsic, the calling convention for any inherited or overriding dispatching operation of a tagged type is that of the corresponding subprogram of the parent type. The default calling convention for a new dispatching operation of a tagged type is the convention of the type.
- 14 Of these four conventions, only Ada and Intrinsic are allowed as a *convention_identifier* in a *pragma Convention*, *Import*, or *Export*.
- 15/2 Two profiles are *type conformant* if they have the same number of parameters, and both have a result if either does, and corresponding parameter and result types are the same, or, for access parameters or access results, corresponding designated types are the same, or corresponding designated profiles are type conformant.
- 16/2 Two profiles are *mode conformant* if they are type-conformant, and corresponding parameters have identical modes, and, for access parameters or access result types, the designated subtypes statically match, or the designated profiles are subtype conformant.
- 17 Two profiles are *subtype conformant* if they are mode-conformant, corresponding subtypes of the profile statically match, and the associated calling conventions are the same. The profile of a generic formal subprogram is not subtype-conformant with any other profile.
- 18 Two profiles are *fully conformant* if they are subtype-conformant, and corresponding parameters have the same names and have *default_expressions* that are fully conformant with one another.
- 19 Two expressions are *fully conformant* if, after replacing each use of an operator with the equivalent *function_call*:
- 20 • each constituent construct of one corresponds to an instance of the same syntactic category in the other, except that an expanded name may correspond to a *direct_name* (or *character_literal*) or to a different expanded name in the other; and
 - 21 • each *direct_name*, *character_literal*, and *selector_name* that is not part of the prefix of an expanded name in one denotes the same declaration as the corresponding *direct_name*, *character_literal*, or *selector_name* in the other; and
- 21.1/1 • each *attribute_designator* in one must be the same as the corresponding *attribute_designator* in the other; and
- 22 • each *primary* that is a literal in one has the same value as the corresponding literal in the other.

Two `known_discriminant_parts` are *fully conformant* if they have the same number of discriminants, and discriminants in the same positions have the same names, statically matching subtypes, and `default_expressions` that are fully conformant with one another.

Two `discrete_subtype_definitions` are *fully conformant* if they are both `subtype_indications` or are both ranges, the `subtype_marks` (if any) denote the same subtype, and the corresponding `simple_expressions` of the ranges (if any) fully conform.

The *prefixed view profile* of a subprogram is the profile obtained by omitting the first parameter of that subprogram. There is no prefixed view profile for a parameterless subprogram. For the purposes of defining subtype and mode conformance, the convention of a prefixed view profile is considered to match that of either an entry or a protected operation.

Implementation Permissions

An implementation may declare an operator declared in a language-defined library unit to be intrinsic.

6.3.2 Inline Expansion of Subprograms

Subprograms may be expanded in line at the call site.

Syntax

The form of a `pragma Inline`, which is a program unit pragma (see 10.1.5), is as follows:

`pragma Inline(name {, name});`

Legality Rules

The `pragma` shall apply to one or more callable entities or generic subprograms.

Static Semantics

If a `pragma Inline` applies to a callable entity, this indicates that inline expansion is desired for all calls to that entity. If a `pragma Inline` applies to a generic subprogram, this indicates that inline expansion is desired for all calls to all instances of that generic subprogram.

Implementation Permissions

For each call, an implementation is free to follow or to ignore the recommendation expressed by the `pragma`.

An implementation may allow a `pragma Inline` that has an argument which is a `direct_name` denoting a `subprogram_body` of the same `declarative_part`.

NOTES

6 The name in a `pragma Inline` can denote more than one entity in the case of overloading. Such a `pragma` applies to all of the denoted entities.

6.4 Subprogram Calls

- 1 A *subprogram call* is either a *procedure_call_statement* or a *function_call*; it invokes the execution of the *subprogram_body*. The call specifies the association of the actual parameters, if any, with formal parameters of the subprogram.

Syntax

- ```

2 procedure_call_statement ::=
3 procedure_name;
| procedure_prefix actual_parameter_part;
4 function_call ::=
5 function_name
| function_prefix actual_parameter_part
6 actual_parameter_part ::=
7 (parameter_association {, parameter_association})
8 parameter_association ::=
9 [formal_parameter_selector_name =>] explicit_actual_parameter
10 explicit_actual_parameter ::= expression | variable_name
11 A parameter_association is named or positional according to whether or not the formal_parameter_selector_name is specified. Any positional associations shall precede any named associations.
12 Named associations are not allowed if the prefix in a subprogram call is an attribute_reference.

```

*Name Resolution Rules*

- 8/2 The name or prefix given in a *procedure\_call\_statement* shall resolve to denote a callable entity that is a procedure, or an entry renamed as (viewed as) a procedure. The name or prefix given in a *function\_call* shall resolve to denote a callable entity that is a function. The name or prefix shall not resolve to denote an abstract subprogram unless it is also a dispatching subprogram. When there is an *actual\_parameter\_part*, the prefix can be an *implicit\_dereference* of an access-to-subprogram value.
- 9 A subprogram call shall contain at most one association for each formal parameter. Each formal parameter without an association shall have a *default\_expression* (in the profile of the view denoted by the name or prefix). This rule is an overloading rule (see 8.6).

*Dynamic Semantics*

- 10/2 For the execution of a subprogram call, the name or prefix of the call is evaluated, and each *parameter\_association* is evaluated (see 6.4.1). If a *default\_expression* is used, an implicit *parameter\_association* is assumed for this rule. These evaluations are done in an arbitrary order. The *subprogram\_body* is then executed, or a call on an entry or protected subprogram is performed (see 3.9.2). Finally, if the subprogram completes normally, then after it is left, any necessary assigning back of formal to actual parameters occurs (see 6.4.1).
- 10.1/2 If the name or prefix of a subprogram call denotes a prefixed view (see 4.1.3), the subprogram call is equivalent to a call on the underlying subprogram, with the first actual parameter being provided by the prefix of the prefixed view (or the Access attribute of this prefix if the first formal parameter is an access parameter), and the remaining actual parameters given by the *actual\_parameter\_part*, if any.
- 11/2 The exception Program\_Error is raised at the point of a *function\_call* if the function completes normally without executing a return statement.

A function\_call denotes a constant, as defined in 6.5; the nominal subtype of the constant is given by the nominal subtype of the function result. 12/2

#### Examples

*Examples of procedure calls:*

|                                                                                                                                                                                                                                |                                                                                       |                                               |
|--------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|---------------------------------------------------------------------------------------|-----------------------------------------------|
| <pre>Traverse_Tree; Print_Header(128, Title, True); Switch(From =&gt; X, To =&gt; Next); Print_Header(128, Header =&gt; Title, Center =&gt; True); Print_Header(Header =&gt; Title, Center =&gt; True, Pages =&gt; 128);</pre> | <small>-- see 6.1<br/>-- see 6.1<br/>-- see 6.1<br/>-- see 6.1<br/>-- see 6.1</small> | <small>13<br/>14<br/>14<br/>15<br/>15</small> |
|--------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|---------------------------------------------------------------------------------------|-----------------------------------------------|

*Examples of function calls:*

|                                                                                                                                                                           |                   |                   |
|---------------------------------------------------------------------------------------------------------------------------------------------------------------------------|-------------------|-------------------|
| <pre>Dot_Product(U, V)      -- see 6.1 and 6.3 Clock                  -- see 9.6 F.all                 -- presuming F is of an access-to-subprogram type — see 3.10</pre> | <small>17</small> | <small>16</small> |
|---------------------------------------------------------------------------------------------------------------------------------------------------------------------------|-------------------|-------------------|

*Examples of procedures with default expressions:*

|                                                                                                                                                                                                                                                                                          |                   |                   |
|------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|-------------------|-------------------|
| <pre>procedure Activate(Process : in Process_Name;                     After   : in Process_Name := No_Process;                     Wait    : in Duration := 0.0;                     Prior   : in Boolean := False);  procedure Pair(Left, Right : in Person_Name := new Person);</pre> | <small>19</small> | <small>18</small> |
|                                                                                                                                                                                                                                                                                          | <small>20</small> |                   |

*Examples of their calls:*

|                                                                                                                                                                                              |                   |                   |
|----------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|-------------------|-------------------|
| <pre>Activate(X); Activate(X, After =&gt; Y); Activate(X, Wait =&gt; 60.0, Prior =&gt; True); Activate(X, Y, 10.0, False);  Pair; Pair(Left =&gt; new Person, Right =&gt; new Person);</pre> | <small>22</small> | <small>21</small> |
|                                                                                                                                                                                              | <small>23</small> |                   |

#### NOTES

If a default\_expression is used for two or more parameters in a multiple parameter\_specification, the default\_expression is evaluated once for each omitted parameter. Hence in the above examples, the two calls of Pair are equivalent. 24

#### Examples

*Examples of overloaded subprograms:*

|                                                                                                                                              |                   |                   |
|----------------------------------------------------------------------------------------------------------------------------------------------|-------------------|-------------------|
| <pre>procedure Put(X : in Integer); procedure Put(X : in String);  procedure Set(Tint   : in Color); procedure Set(Signal : in Light);</pre> | <small>26</small> | <small>25</small> |
|                                                                                                                                              | <small>27</small> |                   |

*Examples of their calls:*

|                                                                                                                                                                                                                              |                   |                   |
|------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|-------------------|-------------------|
| <pre>Put(28); Put("no possible ambiguity here");  Set(Tint  =&gt; Red); Set(Signal =&gt; Red); Set(Color'(Red));  -- Set(Red) would be ambiguous since Red may -- denote a value either of type Color or of type Light</pre> | <small>29</small> | <small>28</small> |
|                                                                                                                                                                                                                              | <small>30</small> |                   |
|                                                                                                                                                                                                                              | <small>31</small> |                   |

## 6.4.1 Parameter Associations

- 1 A parameter association defines the association between an actual parameter and a formal parameter.

### *Name Resolution Rules*

- 2 The *formal\_parameter\_selector\_name* of a *parameter\_association* shall resolve to denote a *parameter\_specification* of the view being called.
- 3 The *actual\_parameter* is either the *explicit\_actual\_parameter* given in a *parameter\_association* for a given formal parameter, or the corresponding *default\_expression* if no *parameter\_association* is given for the formal parameter. The expected type for an actual parameter is the type of the corresponding formal parameter.
- 4 If the mode is **in**, the actual is interpreted as an **expression**; otherwise, the actual is interpreted only as a **name**, if possible.

### *Legality Rules*

- 5 If the mode is **in out** or **out**, the actual shall be a **name** that denotes a variable.
- 6 The type of the actual parameter associated with an access parameter shall be convertible (see 4.6) to its anonymous access type.

### *Dynamic Semantics*

- 7 For the evaluation of a *parameter\_association*:
- 8 • The actual parameter is first evaluated.
  - 9 • For an access parameter, the *access\_definition* is elaborated, which creates the anonymous access type.
  - 10 • For a parameter (of any mode) that is passed by reference (see 6.2), a view conversion of the actual parameter to the nominal subtype of the formal parameter is evaluated, and the formal parameter denotes that conversion.
  - 11 • For an **in** or **in out** parameter that is passed by copy (see 6.2), the formal parameter object is created, and the value of the actual parameter is converted to the nominal subtype of the formal parameter and assigned to the formal.
  - 12 • For an **out** parameter that is passed by copy, the formal parameter object is created, and:
    - 13 • For an access type, the formal parameter is initialized from the value of the actual, without a constraint check;
    - 14 • For a composite type with discriminants or that has implicit initial values for any subcomponents (see 3.3.1), the behavior is as for an **in out** parameter passed by copy.
    - 15 • For any other type, the formal parameter is uninitialized. If composite, a view conversion of the actual parameter to the nominal subtype of the formal is evaluated (which might raise *Constraint\_Error*), and the actual subtype of the formal is that of the view conversion. If elementary, the actual subtype of the formal is given by its nominal subtype.
- 16 A formal parameter of mode **in out** or **out** with discriminants is constrained if either its nominal subtype or the actual parameter is constrained.

After normal completion and leaving of a subprogram, for each **in out** or **out** parameter that is passed by copy, the value of the formal parameter is converted to the subtype of the variable given as the actual parameter and assigned to it. These conversions and assignments occur in an arbitrary order.

17

## 6.5 Return Statements

A **simple\_return\_statement** or **extended\_return\_statement** (collectively called a *return statement*) is used to complete the execution of the innermost enclosing **subprogram\_body**, **entry\_body**, or **accept\_statement**.

1/2

| <i>Syntax</i>                                                                                                                                                                       | 2/2   |
|-------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|-------|
| <b>simple_return_statement</b> ::= <b>return</b> [expression];                                                                                                                      | 2/2   |
| <b>extended_return_statement</b> ::=                                                                                                                                                | 2.1/2 |
| <b>return</b> <b>defining_identifier</b> : [ <b>aliased</b> ] <b>return_subtype_indication</b> [: expression] [ <b>do</b><br>handled_sequence_of_statements<br><b>end return</b> ]; |       |
| <b>return_subtype_indication</b> ::= <b>subtype_indication</b>   <b>access_definition</b>                                                                                           | 2.2/2 |

### *Name Resolution Rules*

The *result subtype* of a function is the subtype denoted by the **subtype\_mark**, or defined by the **access\_definition**, after the reserved word **return** in the profile of the function. The expected type for the expression, if any, of a **simple\_return\_statement** is the result type of the corresponding function. The expected type for the expression of an **extended\_return\_statement** is that of the **return\_subtype\_indication**.

3/2

### *Legality Rules*

A return statement shall be within a callable construct, and it *applies to* the innermost callable construct or **extended\_return\_statement** that contains it. A return statement shall not be within a body that is within the construct to which the return statement applies.

4/2

A function body shall contain at least one return statement that applies to the function body, unless the function contains **code\_statements**. A **simple\_return\_statement** shall include an **expression** if and only if it applies to a function body. An **extended\_return\_statement** shall apply to a function body.

5/2

For an **extended\_return\_statement** that applies to a function body:

5.1/2

- If the result subtype of the function is defined by a **subtype\_mark**, the **return\_subtype\_indication** shall be a **subtype\_indication**. The type of the **subtype\_indication** shall be the result type of the function. If the result subtype of the function is constrained, then the subtype defined by the **subtype\_indication** shall also be constrained and shall statically match this result subtype. If the result subtype of the function is unconstrained, then the subtype defined by the **subtype\_indication** shall be a definite subtype, or there shall be an **expression**.
- If the result subtype of the function is defined by an **access\_definition**, the **return\_subtype\_indication** shall be an **access\_definition**. The subtype defined by the **access\_definition** shall statically match the result subtype of the function. The accessibility level of this anonymous access subtype is that of the result subtype.

5.2/2

5.3/2

For any return statement that applies to a function body:

5.4/2

- 5.5/2 • If the result subtype of the function is limited, then the expression of the return statement (if any) shall be an aggregate, a function call (or equivalent use of an operator), or a qualified\_expression or parenthesized expression whose operand is one of these.
- 5.6/2 • If the result subtype of the function is class-wide, the accessibility level of the type of the expression of the return statement shall not be statically deeper than that of the master that elaborated the function body. If the result subtype has one or more unconstrained access discriminants, the accessibility level of the anonymous access type of each access discriminant, as determined by the expression of the simple\_return\_statement or the return\_subtype\_- indication, shall not be statically deeper than that of the master that elaborated the function body.

*Static Semantics*

- 5.7/2 Within an extended\_return\_statement, the *return object* is declared with the given defining\_identifier, with the nominal subtype defined by the return\_subtype\_indication.

*Dynamic Semantics*

- 5.8/2 For the execution of an extended\_return\_statement, the subtype\_indication or access\_definition is elaborated. This creates the nominal subtype of the return object. If there is an expression, it is evaluated and converted to the nominal subtype (which might raise Constraint\_Error — see 4.6); the return object is created and the converted value is assigned to the return object. Otherwise, the return object is created and initialized by default as for a stand-alone object of its nominal subtype (see 3.3.1). If the nominal subtype is indefinite, the return object is constrained by its initial value.
  - 6/2 For the execution of a simple\_return\_statement, the expression (if any) is first evaluated, converted to the result subtype, and then is assigned to the anonymous *return object*.
  - 7/2 If the return object has any parts that are tasks, the activation of those tasks does not occur until after the function returns (see 9.2).
  - 8/2 If the result type of a function is a specific tagged type, the tag of the return object is that of the result type. If the result type is class-wide, the tag of the return object is that of the value of the expression. A check is made that the accessibility level of the type identified by the tag of the result is not deeper than that of the master that elaborated the function body. If this check fails, Program\_Error is raised.
- Paragraphs 9 through 20 were deleted.*
- 21/2 If the result subtype of a function has one or more unconstrained access discriminants, a check is made that the accessibility level of the anonymous access type of each access discriminant, as determined by the expression or the return\_subtype\_indication of the function, is not deeper than that of the master that elaborated the function body. If this check fails, Program\_Error is raised.
  - 22/2 For the execution of an extended\_return\_statement, the handled\_sequence\_of\_statements is executed. Within this handled\_sequence\_of\_statements, the execution of a simple\_return\_statement that applies to the extended\_return\_statement causes a transfer of control that completes the extended\_return\_statement. Upon completion of a return statement that applies to a callable construct, a transfer of control is performed which completes the execution of the callable construct, and returns to the caller.
  - 23/2 In the case of a function, the function\_call denotes a constant view of the return object.

*Implementation Permissions*

- 24/2 If the result subtype of a function is unconstrained, and a call on the function is used to provide the initial value of an object with a constrained nominal subtype, Constraint\_Error may be raised at the point of the call (after abandoning the execution of the function body) if, while elaborating the return\_subtype\_-

indication or evaluating the expression of a return statement that applies to the function body, it is determined that the value of the result will violate the constraint of the subtype of this object.

#### Examples

*Examples of return statements:*

|                                                                                                                                                                  |                                                                                                                                                                 |                          |
|------------------------------------------------------------------------------------------------------------------------------------------------------------------|-----------------------------------------------------------------------------------------------------------------------------------------------------------------|--------------------------|
| <pre>return;</pre><br><pre>return Key_Value(Last_Index);</pre><br><pre>return Node : Cell do   Node.Value := Result;   Node.Succ := Next_Node; end return;</pre> | -- in a procedure body, entry_body,<br>-- accept_statement, or extended_return_statement<br>-- in a function body<br>-- in a function body, see 3.10.1 for Cell | 25<br>26/2<br>27<br>28/2 |
|------------------------------------------------------------------------------------------------------------------------------------------------------------------|-----------------------------------------------------------------------------------------------------------------------------------------------------------------|--------------------------|

## 6.5.1 Pragma No\_Return

A pragma No\_Return indicates that a procedure cannot return normally; it may propagate an exception or loop forever.

#### Syntax

The form of a pragma No\_Return, which is a representation pragma (see 13.1), is as follows:

2/2

```
pragma No_Return(procedure_local_name{,procedure_local_name});
```

#### Legality Rules

Each *procedure\_local\_name* shall denote one or more procedures or generic procedures; the denoted entities are *non-returning*. The *procedure\_local\_name* shall not denote a null procedure nor an instance of a generic unit.

A return statement shall not apply to a non-returning procedure or generic procedure.

A procedure shall be non-returning if it overrides a dispatching non-returning procedure. In addition to the places where Legality Rules normally apply (see 12.3), this rule applies also in the private part of an instance of a generic unit.

If a renaming-as-body completes a non-returning procedure declaration, then the renamed procedure shall be non-returning.

#### Static Semantics

If a generic procedure is non-returning, then so are its instances. If a procedure declared within a generic unit is non-returning, then so are the corresponding copies of that procedure in instances.

#### Dynamic Semantics

If the body of a non-returning procedure completes normally, Program\_Error is raised at the point of the call.

#### Examples

|                                                                                                                                                                           |  |      |
|---------------------------------------------------------------------------------------------------------------------------------------------------------------------------|--|------|
| <pre>procedure Fail(Msg : String);    -- raises Fatal_Error exception pragma No_Return(Fail);   -- Inform compiler and reader that procedure never returns normally</pre> |  | 10/2 |
|---------------------------------------------------------------------------------------------------------------------------------------------------------------------------|--|------|

## 6.6 Overloading of Operators

- 1 An *operator* is a function whose designator is an `operator_symbol`. Operators, like other functions, may be overloaded.

### *Name Resolution Rules*

- 2 Each use of a unary or binary operator is equivalent to a `function_call` with `function_prefix` being the corresponding `operator_symbol`, and with (respectively) one or two positional actual parameters being the operand(s) of the operator (in order).

### *Legality Rules*

- 3 The `subprogram_specification` of a unary or binary operator shall have one or two parameters, respectively. A generic function instantiation whose `designator` is an `operator_symbol` is only allowed if the specification of the generic function has the corresponding number of parameters.
- 4 `Default_expressions` are not allowed for the parameters of an operator (whether the operator is declared with an explicit `subprogram_specification` or by a `generic_instantiation`).
- 5 An explicit declaration of `"=/"` shall not have a result type of the predefined type `Boolean`.

### *Static Semantics*

- 6 A declaration of `"=`" whose result type is `Boolean` implicitly declares a declaration of `"=/"` that gives the complementary result.

#### NOTES

- 7 8 The operators `"+"` and `"-"` are both unary and binary operators, and hence may be overloaded with both one- and two-parameter functions.

### *Examples*

- 8 Examples of user-defined operators:

```
9 function "+" (Left, Right : Matrix) return Matrix;
 function "+" (Left, Right : Vector) return Vector;
```

-- assuming that *A*, *B*, and *C* are of the type `Vector`  
-- the following two statements are equivalent:

```
A := B + C;
A := "+"(B, C);
```

## 6.7 Null Procedures

A `null_procedure_declaration` provides a shorthand to declare a procedure with an empty body.

1/2

### Syntax

```
null_procedure_declaration ::=
 [overriding_indicator]
 procedure_specification is null;
```

2/2

### Static Semantics

A `null_procedure_declaration` declares a *null procedure*. A completion is not allowed for a `null_procedure_declaration`.

3/2

### Dynamic Semantics

The execution of a *null procedure* is invoked by a subprogram call. For the execution of a subprogram call on a *null procedure*, the execution of the `subprogram_body` has no effect.

4/2

The elaboration of a `null_procedure_declaration` has no effect.

5/2

### Examples

```
procedure Simplify(Expr : in out Expression) is null; -- see 3.9
-- By default, Simplify does nothing, but it may be overridden in extensions of Expression
```

6/2

## Section 7: Packages

Packages are program units that allow the specification of groups of logically related entities. Typically, a package contains the declaration of a type (often a private type or private extension) along with the declarations of primitive subprograms of the type, which can be called from outside the package, while their inner workings remain hidden from outside users.

### 7.1 Package Specifications and Declarations

A package is generally provided in two parts: a `package_specification` and a `package_body`. Every package has a `package_specification`, but not all packages have a `package_body`.

*Syntax*

```
package_declaration ::= package_specification;
package_specification ::=
 package defining_program_unit_name is
 {basic_declarative_item}
 [private
 {basic_declarative_item}]
 end [[parent_unit_name.]identifier]
```

If an identifier or `parent_unit_name.identifier` appears at the end of a `package_specification`, then this sequence of lexical elements shall repeat the `defining_program_unit_name`.

*Legality Rules*

A `package_declaration` or `generic_package_declaration` requires a completion (a body) if it contains any `basic_declarative_item` that requires a completion, but whose completion is not in its `package_specification`.

*Static Semantics*

The first list of `basic_declarative_items` of a `package_specification` of a package other than a generic formal package is called the *visible part* of the package. The optional list of `basic_declarative_items` after the reserved word `private` (of any `package_specification`) is called the *private part* of the package. If the reserved word `private` does not appear, the package has an implicit empty private part. Each list of `basic_declarative_items` of a `package_specification` forms a *declaration list* of the package.

An entity declared in the private part of a package is visible only within the declarative region of the package itself (including any child units — see 10.1.1). In contrast, expanded names denoting entities declared in the visible part can be used even outside the package; furthermore, direct visibility of such entities can be achieved by means of `use_clauses` (see 4.1.3 and 8.4).

*Dynamic Semantics*

The elaboration of a `package_declaration` consists of the elaboration of its `basic_declarative_items` in the given order.

**NOTES**

1 The visible part of a package contains all the information that another program unit is able to know about the package.

2 If a declaration occurs immediately within the specification of a package, and the declaration has a corresponding completion that is a body, then that body has to occur immediately within the body of the package.

*Examples*

11    Example of a package declaration:

```

12 package Rational_Numbers is
13 type Rational is
14 record
15 Numerator : Integer;
16 Denominator : Positive;
17 end record;
18
19 function "=" (X, Y : Rational) return Boolean;
20 function "/" (X, Y : Integer) return Rational; -- to construct a rational number
21 function "+" (X, Y : Rational) return Rational;
22 function "-" (X, Y : Rational) return Rational;
23 function "*" (X, Y : Rational) return Rational;
24 function "/" (X, Y : Rational) return Rational;
25
26 end Rational_Numbers;
```

17    There are also many examples of package declarations in the predefined language environment (see Annex A).

## 7.2 Package Bodies

1    In contrast to the entities declared in the visible part of a package, the entities declared in the package\_body are visible only within the package\_body itself. As a consequence, a package with a package\_body can be used for the construction of a group of related subprograms in which the logical operations available to clients are clearly isolated from the internal entities.

2    *Syntax*

```

package_body ::=
 package body defining_program_unit_name is
 declarative_part
 [begin
 handled_sequence_of_statements]
 end [[parent_unit_name.]identifier];
```

3    If an identifier or parent\_unit\_name.identifier appears at the end of a package\_body, then this sequence of lexical elements shall repeat the defining\_program\_unit\_name.

*Legality Rules*

4    A package\_body shall be the completion of a previous package\_declaration or generic\_package\_declaration. A library package\_declaration or library generic\_package\_declaration shall not have a body unless it requires a body; **pragma Elaborate\_Body** can be used to require a library\_unit\_declaration to have a body (see 10.2.1) if it would not otherwise require one.

*Static Semantics*

5    In any package\_body without statements there is an implicit null\_statement. For any package\_declaration without an explicit completion, there is an implicit package\_body containing a single null\_statement. For a noninstance, nonlibrary package, this body occurs at the end of the declarative\_part of the innermost enclosing program unit or block\_statement; if there are several such packages, the order of the implicit package\_bodies is unspecified. (For an instance, the implicit package\_body occurs at the place of the instantiation (see 12.3). For a library package, the place is partially determined by the elaboration dependences (see Section 10).)

*Dynamic Semantics*

For the elaboration of a nongeneric `package_body`, its `declarative_part` is first elaborated, and its `handled_sequence_of_statements` is then executed.

## NOTES

3 A variable declared in the body of a package is only visible within this body and, consequently, its value can only be changed within the `package_body`. In the absence of local tasks, the value of such a variable remains unchanged between calls issued from outside the package to subprograms declared in the visible part. The properties of such a variable are similar to those of a “static” variable of C.

4 The elaboration of the body of a subprogram explicitly declared in the visible part of a package is caused by the elaboration of the body of the package. Hence a call of such a subprogram by an outside program unit raises the exception `Program_Error` if the call takes place before the elaboration of the `package_body` (see 3.11).

*Examples*

*Example of a package body (see 7.1):*

```

package body Rational_Numbers is
 procedure Same_Denominator (X,Y : in out Rational) is
 begin
 -- reduces X and Y to the same denominator:
 ...
 end Same_Denominator;
 function "="(X,Y : Rational) return Boolean is
 U : Rational := X;
 V : Rational := Y;
 begin
 Same_Denominator (U,V);
 return U.Numerator = V.Numerator;
 end "=";
 function "/" (X,Y : Integer) return Rational is
 begin
 if Y > 0 then
 return (Numerator => X, Denominator => Y);
 else
 return (Numerator => -X, Denominator => -Y);
 end if;
 end "/";
 function "+" (X,Y : Rational) return Rational is ... end "+";
 function "-" (X,Y : Rational) return Rational is ... end "-";
 function "*" (X,Y : Rational) return Rational is ... end "*";
 function "/" (X,Y : Rational) return Rational is ... end "/";
end Rational_Numbers;
```

## 7.3 Private Types and Private Extensions

The declaration (in the visible part of a package) of a type as a private type or private extension serves to separate the characteristics that can be used directly by outside program units (that is, the logical properties) from other characteristics whose direct use is confined to the package (the details of the definition of the type itself). See 3.9.1 for an overview of type extensions.

*Syntax*

```

private_type_declaration ::=
 type defining_identifier [discriminant_part] is [[abstract] tagged] [limited] private;
```

3/2      private\_extension\_declaration ::=  
           type defining\_identifier [discriminant\_part] is  
           [abstract] [limited | synchronized] new ancestor\_subtype\_indication  
           [and interface\_list] with private;

*Legality Rules*

- 4 A private\_type\_declaration or private\_extension\_declaration declares a *partial view* of the type; such a declaration is allowed only as a declarative\_item of the visible part of a package, and it requires a completion, which shall be a full\_type\_declaration that occurs as a declarative\_item of the private part of the package. The view of the type declared by the full\_type\_declaration is called the *full view*. A generic formal private type or a generic formal private extension is also a partial view.
- 5 A type shall be completely defined before it is frozen (see 3.11.1 and 13.14). Thus, neither the declaration of a variable of a partial view of a type, nor the creation by an allocator of an object of the partial view are allowed before the full declaration of the type. Similarly, before the full declaration, the name of the partial view cannot be used in a generic\_instantiation or in a representation item.
- 6/2 A private type is limited if its declaration includes the reserved word **limited**; a private extension is limited if its ancestor type is a limited type that is not an interface type, or if the reserved word **limited** or **synchronized** appears in its definition. If the partial view is nonlimited, then the full view shall be nonlimited. If a tagged partial view is limited, then the full view shall be limited. On the other hand, if an untagged partial view is limited, the full view may be limited or nonlimited.
- 7 If the partial view is tagged, then the full view shall be tagged. On the other hand, if the partial view is untagged, then the full view may be tagged or untagged. In the case where the partial view is untagged and the full view is tagged, no derivatives of the partial view are allowed within the immediate scope of the partial view; derivatives of the full view are allowed.
- 7.1/2 If a full type has a partial view that is tagged, then:
  - 7.2/2 • the partial view shall be a synchronized tagged type (see 3.9.4) if and only if the full type is a synchronized tagged type;
  - 7.3/2 • the partial view shall be a descendant of an interface type (see 3.9.4) if and only if the full type is a descendant of the interface type.
- 8 The *ancestor\_subtype* of a private\_extension\_declaration is the subtype defined by the *ancestor\_subtype\_indication*; the ancestor type shall be a specific tagged type. The full view of a private extension shall be derived (directly or indirectly) from the ancestor type. In addition to the places where Legality Rules normally apply (see 12.3), the requirement that the ancestor be specific applies also in the private part of an instance of a generic unit.
- 8.1/2 If the reserved word **limited** appears in a private\_extension\_declaration, the ancestor type shall be a limited type. If the reserved word **synchronized** appears in a private\_extension\_declaration, the ancestor type shall be a limited interface.
- 9 If the declaration of a partial view includes a known\_discriminant\_part, then the full\_type\_declaration shall have a fully conforming (explicit) known\_discriminant\_part (see 6.3.1, “Conformance Rules”). The ancestor subtype may be unconstrained; the parent subtype of the full view is required to be constrained (see 3.7).
- 10 If a private extension inherits known discriminants from the ancestor subtype, then the full view shall also inherit its discriminants from the ancestor subtype, and the parent subtype of the full view shall be constrained if and only if the ancestor subtype is constrained.

If the `full_type_declaration` for a private extension is defined by a `derived_type_definition`, then the reserved word `limited` shall appear in the `full_type_declaration` if and only if it also appears in the `private_extension_declaration`. 10.1/2

If a partial view has unknown discriminants, then the `full_type_declaration` may define a definite or an indefinite subtype, with or without discriminants. 11

If a partial view has neither known nor unknown discriminants, then the `full_type_declaration` shall define a definite subtype. 12

If the ancestor subtype of a private extension has constrained discriminants, then the parent subtype of the full view shall impose a statically matching constraint on those discriminants. 13

#### *Static Semantics*

A `private_type_declaration` declares a private type and its first subtype. Similarly, a `private_extension_declaration` declares a private extension and its first subtype. 14

A declaration of a partial view and the corresponding `full_type_declaration` define two views of a single type. The declaration of a partial view together with the visible part define the operations that are available to outside program units; the declaration of the full view together with the private part define other operations whose direct use is possible only within the declarative region of the package itself. Moreover, within the scope of the declaration of the full view, the *characteristics* of the type are determined by the full view; in particular, within its scope, the full view determines the classes that include the type, which components, entries, and protected subprograms are visible, what attributes and other predefined operations are allowed, and whether the first subtype is static. See 7.3.1. 15

A private extension inherits components (including discriminants unless there is a new `discriminant_part` specified) and user-defined primitive subprograms from its ancestor type and its progenitor types (if any), in the same way that a record extension inherits components and user-defined primitive subprograms from its parent type and its progenitor types (see 3.4). 16/2

#### *Dynamic Semantics*

The elaboration of a `private_type_declaration` creates a partial view of a type. The elaboration of a `private_extension_declaration` elaborates the *ancestor\_subtype\_indication*, and creates a partial view of a type. 17

#### NOTES

5 The partial view of a type as declared by a `private_type_declaration` is defined to be a composite view (in 3.2). The full view of the type might or might not be composite. A private extension is also composite, as is its full view. 18

6 Declaring a private type with an `unknown_discriminant_part` is a way of preventing clients from creating uninitialized objects of the type; they are then forced to initialize each object by calling some operation declared in the visible part of the package. 19/2

7 The ancestor type specified in a `private_extension_declaration` and the parent type specified in the corresponding declaration of a record extension given in the private part need not be the same. If the ancestor type is not an interface type, the parent type of the full view can be any descendant of the ancestor type. In this case, for a primitive subprogram that is inherited from the ancestor type and not overridden, the formal parameter names and default expressions (if any) come from the corresponding primitive subprogram of the specified ancestor type, while the body comes from the corresponding primitive subprogram of the parent type of the full view. See 3.9.2. 20/2

8 If the ancestor type specified in a `private_extension_declaration` is an interface type, the parent type can be any type so long as the full view is a descendant of the ancestor type. The progenitor types specified in a `private_extension_declaration` and the progenitor types specified in the corresponding declaration of a record extension given in the private part need not be the same — the only requirement is that the private extension and the record extension be descended from the same set of interfaces. 20.1/2

*Examples*

21 Examples of private type declarations:

22    `type Key is private;`  
      `type File_Name is limited private;`

23 Example of a private extension declaration:

24    `type List is new Ada.Finalization.Controlled with private;`

### 7.3.1 Private Operations

- 1 For a type declared in the visible part of a package or generic package, certain operations on the type do not become visible until later in the package — either in the private part or the body. Such *private operations* are available only inside the declarative region of the package or generic package.

*Static Semantics*

- 2 The predefined operators that exist for a given type are determined by the classes to which the type belongs. For example, an integer type has a predefined "+" operator. In most cases, the predefined operators of a type are declared immediately after the definition of the type; the exceptions are explained below. Inherited subprograms are also implicitly declared immediately after the definition of the type, except as stated below.

- 3/1 For a composite type, the characteristics (see 7.3) of the type are determined in part by the characteristics of its component types. At the place where the composite type is declared, the only characteristics of component types used are those characteristics visible at that place. If later immediately within the declarative region in which the composite type is declared additional characteristics become visible for a component type, then any corresponding characteristics become visible for the composite type. Any additional predefined operators are implicitly declared at that place.

- 4/1 The corresponding rule applies to a type defined by a `derived_type_definition`, if there is a place immediately within the declarative region in which the type is declared where additional characteristics of its parent type become visible.

- 5/1 For example, an array type whose component type is limited private becomes nonlimited if the full view of the component type is nonlimited and visible at some later place immediately within the declarative region in which the array type is declared. In such a case, the predefined "=" operator is implicitly declared at that place, and assignment is allowed after that place.

- 6/1 Inherited primitive subprograms follow a different rule. For a `derived_type_definition`, each inherited primitive subprogram is implicitly declared at the earliest place, if any, immediately within the declarative region in which the `type_declaration` occurs, but after the `type_declaration`, where the corresponding declaration from the parent is visible. If there is no such place, then the inherited subprogram is not declared at all. An inherited subprogram that is not declared at all cannot be named in a call and cannot be overridden, but for a tagged type, it is possible to dispatch to it.

- 7 For a `private_extension_declaration`, each inherited subprogram is declared immediately after the `private_extension_declaration` if the corresponding declaration from the ancestor is visible at that place. Otherwise, the inherited subprogram is not declared for the private extension, though it might be for the full type.

- 8 The `Class` attribute is defined for tagged subtypes in 3.9. In addition, for every subtype S of an untagged private type whose full view is tagged, the following attribute is defined:

|         |                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                |    |
|---------|----------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|----|
| S'Class | Denotes the class-wide subtype corresponding to the full view of S. This attribute is allowed only from the beginning of the private part in which the full view is declared, until the declaration of the full view. After the full view, the Class attribute of the full view can be used.                                                                                                                                                                                                                                                                                                                   | 9  |
| NOTES   |                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                |    |
| 9       | Because a partial view and a full view are two different views of one and the same type, outside of the defining package the characteristics of the type are those defined by the visible part. Within these outside program units the type is just a private type or private extension, and any language rule that applies only to another class of types does not apply. The fact that the full declaration might implement a private type with a type of a particular class (for example, as an array type) is relevant only within the declarative region of the package itself including any child units. | 10 |
| 11      | The consequences of this actual implementation are, however, valid everywhere. For example: any default initialization of components takes place; the attribute Size provides the size of the full view; finalization is still done for controlled components of the full view; task dependence rules still apply to components that are task objects.                                                                                                                                                                                                                                                         |    |
| 12/2    | 10 Partial views provide initialization, membership tests, selected components for the selection of discriminants and inherited components, qualification, and explicit conversion. Nonlimited partial views also allow use of assignment_statements.                                                                                                                                                                                                                                                                                                                                                          |    |
| 13      | 11 For a subtype S of a partial view, S'Size is defined (see 13.3). For an object A of a partial view, the attributes A'Size and A'Address are defined (see 13.3). The Position, First_Bit, and Last_Bit attributes are also defined for discriminants and inherited components.                                                                                                                                                                                                                                                                                                                               |    |
| 14      | <i>Examples</i>                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                |    |
|         | <i>Example of a type with private operations:</i>                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                              |    |
| 15      | <pre>package Key_Manager is     type Key is private;     Null_Key : constant Key; -- a deferred constant declaration (see 7.4)     procedure Get_Key(K : out Key);     function "&lt;" (X, Y : Key) return Boolean; private     type Key is new Natural;     Null_Key : constant Key := Key'First; end Key_Manager;</pre>                                                                                                                                                                                                                                                                                      |    |
| 16      | <pre>package body Key_Manager is     Last_Key : Key := Null_Key;     procedure Get_Key(K : out Key) is     begin         Last_Key := Last_Key + 1;         K := Last_Key;     end Get_Key;     function "&lt;" (X, Y : Key) return Boolean is     begin         return Natural(X) &lt; Natural(Y);     end "&lt;"; end Key_Manager;</pre>                                                                                                                                                                                                                                                                      | 17 |
| 18      | NOTES                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                          |    |
| 12      | <i>Notes on the example.</i> Outside of the package Key_Manager, the operations available for objects of type Key include assignment, the comparison for equality or inequality, the procedure Get_Key and the operator "<", they do not include other relational operators such as ">=", or arithmetic operators.                                                                                                                                                                                                                                                                                             |    |
| 19      | The explicitly declared operator "<" hides the predefined operator "<" implicitly declared by the full_type_declaration. Within the body of the function, an explicit conversion of X and Y to the subtype Natural is necessary to invoke the "<" operator of the parent type. Alternatively, the result of the function could be written as not (X >= Y), since the operator ">=" is not redefined.                                                                                                                                                                                                           |    |
| 20      | The value of the variable Last_Key, declared in the package body, remains unchanged between calls of the procedure Get_Key. (See also the NOTES of 7.2.)                                                                                                                                                                                                                                                                                                                                                                                                                                                       |    |

## 7.4 Deferred Constants

- 1 Deferred constant declarations may be used to declare constants in the visible part of a package, but with the value of the constant given in the private part. They may also be used to declare constants imported from other languages (see Annex B).

### *Legality Rules*

- 2 A *deferred constant declaration* is an *object\_declaration* with the reserved word **constant** but no initialization expression. The constant declared by a deferred constant declaration is called a *deferred constant*. A deferred constant declaration requires a completion, which shall be a full constant declaration (called the *full declaration* of the deferred constant), or a **pragma Import** (see Annex B).
- 3 A deferred constant declaration that is completed by a full constant declaration shall occur immediately within the visible part of a *package\_specification*. For this case, the following additional rules apply to the corresponding full declaration:
- 4 • The full declaration shall occur immediately within the private part of the same package;
  - 5/2 • The deferred and full constants shall have the same type, or shall have statically matching anonymous access subtypes;
  - 6/2 • If the deferred constant declaration includes a *subtype\_indication* that defines a constrained subtype, then the subtype defined by the *subtype\_indication* in the full declaration shall match it statically. On the other hand, if the subtype of the deferred constant is unconstrained, then the full declaration is still allowed to impose a constraint. The constant itself will be constrained, like all constants;
  - 7/2 • If the deferred constant declaration includes the reserved word **aliased**, then the full declaration shall also;
  - 7.1/2 • If the subtype of the deferred constant declaration excludes null, the subtype of the full declaration shall also exclude null.
- 8 A deferred constant declaration that is completed by a **pragma Import** need not appear in the visible part of a *package\_specification*, and has no full constant declaration.
- 9/2 The completion of a deferred constant declaration shall occur before the constant is frozen (see 13.14).

### *Dynamic Semantics*

- 10 The elaboration of a deferred constant declaration elaborates the *subtype\_indication* or (only allowed in the case of an imported constant) the *array\_type\_definition*.

#### NOTES

- 11 13 The full constant declaration for a deferred constant that is of a given private type or private extension is not allowed before the corresponding *full\_type\_declaration*. This is a consequence of the freezing rules for types (see 13.14).

### *Examples*

- 12 Examples of deferred constant declarations:

```
13 Null_Key : constant Key; -- see 7.3.1
14 CPU_Identifier : constant String(1..8);
 pragma Import(Assembler, CPU_Identifier, Link_Name => "CPU_ID");
 -- see B.1
```

## 7.5 Limited Types

A limited type is (a view of) a type for which copying (such as for an `assignment_statement`) is not allowed. A nonlimited type is a (view of a) type for which copying is allowed.

### *Legality Rules*

If a tagged record type has any limited components, then the reserved word **limited** shall appear in its `record_type_definition`. If the reserved word **limited** appears in the definition of a `derived_type_definition`, its parent type and any progenitor interfaces shall be limited.

In the following contexts, an expression of a limited type is not permitted unless it is an `aggregate`, a `function_call`, or a parenthesized `expression` or `qualified_expression` whose operand is permitted by this rule:

- the initialization `expression` of an `object_declaration` (see 3.3.1) 2.2/2
- the `default_expression` of a `component_declaration` (see 3.8) 2.3/2
- the `expression` of a `record_component_association` (see 4.3.1) 2.4/2
- the `expression` for an `ancestor_part` of an `extension_aggregate` (see 4.3.2) 2.5/2
- an `expression` of a `positional_array_aggregate` or the `expression` of an `array_component_association` (see 4.3.3) 2.6/2
- the `qualified_expression` of an initialized allocator (see 4.8) 2.7/2
- the `expression` of a return statement (see 6.5) 2.8/2
- the `default_expression` or actual parameter for a formal object of mode **in** (see 12.4) 2.9/2

### *Static Semantics*

A type is *limited* if it is one of the following:

- a type with the reserved word **limited**, **synchronized**, **task**, or **protected** in its definition; 4/2
- *This paragraph was deleted.* 5/2
- a composite type with a limited component; 6/2
- a derived type whose parent is limited and is not an interface. 6.1/2

Otherwise, the type is nonlimited.

There are no predefined equality operators for a limited type.

### *Implementation Requirements*

For an `aggregate` of a limited type used to initialize an object as allowed above, the implementation shall not create a separate anonymous object for the `aggregate`. For a `function_call` of a type with a part that is of a task, protected, or explicitly limited record type that is used to initialize an object as allowed above, the implementation shall not create a separate return object (see 6.5) for the `function_call`. The `aggregate` or `function_call` shall be constructed directly in the new object.

#### NOTES

14 While it is allowed to write initializations of limited objects, such initializations never copy a limited object. The source of such an assignment operation must be an `aggregate` or `function_call`, and such `aggregates` and `function_calls` must be built directly in the target object.

*Paragraphs 10 through 15 were deleted.*

16        15 As illustrated in 7.3.1, an untagged limited type can become nonlimited under certain circumstances.

*Examples*

17        Example of a package with a limited type:

```

18 package IO_Package is
19 type File_Name is limited private;
20 procedure Open (F : in out File_Name);
21 procedure Close(F : in out File_Name);
22 procedure Read (F : in File_Name; Item : out Integer);
23 procedure Write(F : in File_Name; Item : in Integer);
24 private
25 type File_Name is
26 limited record
27 Internal_Name : Integer := 0;
28 end record;
29 end IO_Package;
30
31 package body IO_Package is
32 Limit : constant := 200;
33 type File_Descriptor is record ... end record;
34 Directory : array (1 .. Limit) of File_Descriptor;
35 ...
36 procedure Open (F : in out File_Name) is ... end;
37 procedure Close(F : in out File_Name) is ... end;
38 procedure Read (F : in File_Name; Item : out Integer) is ... end;
39 procedure Write(F : in File_Name; Item : in Integer) is ... end;
40 begin
41 ...
42 end IO_Package;

```

NOTES

21        16 *Notes on the example:* In the example above, an outside subprogram making use of IO\_Package may obtain a file name by calling Open and later use it in calls to Read and Write. Thus, outside the package, a file name obtained from Open acts as a kind of password; its internal properties (such as containing a numeric value) are not known and no other operations (such as addition or comparison of internal names) can be performed on a file name. Most importantly, clients of the package cannot make copies of objects of type File\_Name.

22        This example is characteristic of any case where complete control over the operations of a type is desired. Such packages serve a dual purpose. They prevent a user from making use of the internal structure of the type. They also implement the notion of an encapsulated data type where the only operations on the type are those given in the package specification.

23/2      The fact that the full view of File\_Name is explicitly declared **limited** means that parameter passing will always be by reference and function results will always be built directly in the result object (see 6.2 and 6.5).

## 7.6 User-Defined Assignment and Finalization

1        Three kinds of actions are fundamental to the manipulation of objects: initialization, finalization, and assignment. Every object is initialized, either explicitly or by default, after being created (for example, by an **object\_declaration** or **allocator**). Every object is finalized before being destroyed (for example, by leaving a **subprogram\_body** containing an **object\_declaration**, or by a call to an instance of **Unchecked\_Deallocation**). An assignment operation is used as part of **assignment\_statements**, explicit initialization, parameter passing, and other operations.

2        Default definitions for these three fundamental operations are provided by the language, but a *controlled* type gives the user additional control over parts of these operations. In particular, the user can define, for a controlled type, an Initialize procedure which is invoked immediately after the normal default initialization of a controlled object, a Finalize procedure which is invoked immediately before finalization of any of the components of a controlled object, and an Adjust procedure which is invoked as the last step of an assignment to a (nonlimited) controlled object.

*Static Semantics*

The following language-defined library package exists:

```

package Ada.Finalization is
 pragma Preelaborate(Finalization);
 pragma Remote_Types(Finalization);

 type Controlled is abstract tagged private;
 pragma Preelaborable_Initialization(Controlled);

 procedure Initialize (Object : in out Controlled) is null;
 procedure Adjust (Object : in out Controlled) is null;
 procedure Finalize (Object : in out Controlled) is null;

 type Limited_Controlled is abstract tagged limited private;
 pragma Preelaborable_Initialization(Limited_Controlled);

 procedure Initialize (Object : in out Limited_Controlled) is null;
 procedure Finalize (Object : in out Limited_Controlled) is null;
private
 ... -- not specified by the language
end Ada.Finalization;
```

A controlled type is a descendant of Controlled or Limited\_Controlled. The predefined "`=`" operator of type Controlled always returns True, since this operator is incorporated into the implementation of the predefined equality operator of types derived from Controlled, as explained in 4.5.2. The type Limited\_Controlled is like Controlled, except that it is limited and it lacks the primitive subprogram Adjust.

A type is said to *need finalization* if:

- it is a controlled type, a task type or a protected type; or
- it has a component that needs finalization; or
- it is a limited type that has an access discriminant whose designated type needs finalization; or
- it is one of a number of language-defined types that are explicitly defined to need finalization.

*Dynamic Semantics*

During the elaboration or evaluation of a construct that causes an object to be initialized by default, for every controlled subcomponent of the object that is not assigned an initial value (as defined in 3.3.1), Initialize is called on that subcomponent. Similarly, if the object that is initialized by default as a whole is controlled, Initialize is called on the object.

For an extension\_aggregate whose ancestor\_part is a subtype\_mark denoting a controlled subtype, the Initialize procedure of the ancestor type is called, unless that Initialize procedure is abstract.

Initialize and other initialization operations are done in an arbitrary order, except as follows. Initialize is applied to an object after initialization of its subcomponents, if any (including both implicit initialization and Initialize calls). If an object has a component with an access discriminant constrained by a per-object expression, Initialize is applied to this component after any components that do not have such discriminants. For an object with several components with such a discriminant, Initialize is applied to them in order of their component\_declarations. For an allocator, any task activations follow all calls on Initialize.

When a target object with any controlled parts is assigned a value, either when created or in a subsequent assignment\_statement, the *assignment operation* proceeds as follows:

- The value of the target becomes the assigned value.
- The value of the target is *adjusted*.

- 16 To adjust the value of a (nonlimited) composite object, the values of the components of the object are first adjusted in an arbitrary order, and then, if the object is controlled, Adjust is called. Adjusting the value of an elementary object has no effect, nor does adjusting the value of a composite object with no controlled parts.
- 17 For an **assignment\_statement**, after the **name** and **expression** have been evaluated, and any conversion (including constraint checking) has been done, an anonymous object is created, and the value is assigned into it; that is, the assignment operation is applied. (Assignment includes value adjustment.) The target of the **assignment\_statement** is then finalized. The value of the anonymous object is then assigned into the target of the **assignment\_statement**. Finally, the anonymous object is finalized. As explained below, the implementation may eliminate the intermediate anonymous object, so this description subsumes the one given in 5.2, “Assignment Statements”.

*Implementation Requirements*

- 17.1/2 For an **aggregate** of a controlled type whose value is assigned, other than by an **assignment\_statement**, the implementation shall not create a separate anonymous object for the **aggregate**. The aggregate value shall be constructed directly in the target of the assignment operation and Adjust is not called on the target object.

*Implementation Permissions*

- 18 An implementation is allowed to relax the above rules (for nonlimited controlled types) in the following ways:
  - 19 • For an **assignment\_statement** that assigns to an object the value of that same object, the implementation need not do anything.
  - 20 • For an **assignment\_statement** for a noncontrolled type, the implementation may finalize and assign each component of the variable separately (rather than finalizing the entire variable and assigning the entire new value) unless a discriminant of the variable is changed by the assignment.
  - 21/2 • For an **aggregate** or function call whose value is assigned into a target object, the implementation need not create a separate anonymous object if it can safely create the value of the **aggregate** or function call directly in the target object. Similarly, for an **assignment\_statement**, the implementation need not create an anonymous object if the value being assigned is the result of evaluating a **name** denoting an object (the source object) whose storage cannot overlap with the target. If the source object might overlap with the target object, then the implementation can avoid the need for an intermediary anonymous object by exercising one of the above permissions and perform the assignment one component at a time (for an overlapping array assignment), or not at all (for an assignment where the target and the source of the assignment are the same object).
- 22/2 Furthermore, an implementation is permitted to omit implicit Initialize, Adjust, and Finalize calls and associated assignment operations on an object of a nonlimited controlled type provided that:
  - 23/2 • any omitted Initialize call is not a call on a user-defined Initialize procedure, and
  - 24/2 • any usage of the value of the object after the implicit Initialize or Adjust call and before any subsequent Finalize call on the object does not change the external effect of the program, and
  - 25/2 • after the omission of such calls and operations, any execution of the program that executes an Initialize or Adjust call on an object or initializes an object by an **aggregate** will also later execute a Finalize call on the object and will always do so prior to assigning a new value to the object, and
  - 26/2 • the assignment operations associated with omitted Adjust calls are also omitted.

This permission applies to Adjust and Finalize calls even if the implicit calls have additional external effects. 27/2

## 7.6.1 Completion and Finalization

This subclause defines *completion* and *leaving* of the execution of constructs and entities. A *master* is the execution of a construct that includes finalization of local objects after it is complete (and after waiting for any local tasks — see 9.3), but before leaving. Other constructs and entities are left immediately upon completion.

### *Dynamic Semantics*

The execution of a construct or entity is *complete* when the end of that execution has been reached, or when a transfer of control (see 5.1) causes it to be abandoned. Completion due to reaching the end of execution, or due to the transfer of control of an `exit_statement`, `return statement`, `goto_statement`, or `requeue_statement` or of the selection of a `terminate_alternative` is *normal completion*. Completion is *abnormal* otherwise — when control is transferred out of a construct due to abort or the raising of an exception. 2/2

After execution of a construct or entity is complete, it is *left*, meaning that execution continues with the next action, as defined for the execution that is taking place. Leaving an execution happens immediately after its completion, except in the case of a *master*: the execution of a body other than a `package_body`; the execution of a `statement`; or the evaluation of an `expression`, `function_call`, or `range` that is not part of an enclosing `expression`, `function_call`, `range`, or `simple_statement` other than a `simple_return_statement`. A master is finalized after it is complete, and before it is left. 3/2

For the *finalization* of a master, dependent tasks are first awaited, as explained in 9.3. Then each object whose accessibility level is the same as that of the master is finalized if the object was successfully initialized and still exists. These actions are performed whether the master is left by reaching the last statement or via a transfer of control. When a transfer of control causes completion of an execution, each included master is finalized in order, from innermost outward. 4

For the *finalization* of an object:

- If the object is of an elementary type, finalization has no effect; 5
- If the object is of a controlled type, the `Finalize` procedure is called; 6
- If the object is of a protected type, the actions defined in 9.4 are performed; 7
- If the object is of a composite type, then after performing the above actions, if any, every component of the object is finalized in an arbitrary order, except as follows: if the object has a component with an access discriminant constrained by a per-object expression, this component is finalized before any components that do not have such discriminants; for an object with several components with such a discriminant, they are finalized in the reverse of the order of their `component_declarations`; 8
- If the object has coextensions (see 3.10.2), each coextension is finalized after the object whose access discriminant designates it. 9/2

Immediately before an instance of `Unchecked_Deallocation` reclaims the storage of an object, the object is finalized. If an instance of `Unchecked_Deallocation` is never applied to an object created by an allocator, the object will still exist when the corresponding master completes, and it will be finalized then. 10

The order in which the finalization of a master performs finalization of objects is as follows: Objects created by declarations in the master are finalized in the reverse order of their creation. For objects that

were created by allocators for an access type whose ultimate ancestor is declared in the master, this rule is applied as though each such object that still exists had been created in an arbitrary order at the first freezing point (see 13.14) of the ultimate ancestor type; the finalization of these objects is called the *finalization of the collection*. After the finalization of a master is complete, the objects finalized as part of its finalization cease to exist, as do any types and subtypes defined and created within the master.

- 12/2 The target of an **assignment\_statement** is finalized before copying in the new value, as explained in 7.6.
- 13/2 The master of an object is the master enclosing its creation whose accessibility level (see 3.10.2) is equal to that of the object.
- 13.1/2 In the case of an **expression** that is a master, finalization of any (anonymous) objects occurs as the final part of evaluation of the **expression**.

#### *Bounded (Run-Time) Errors*

- 14/1 It is a bounded error for a call on **Finalize** or **Adjust** that occurs as part of object finalization or assignment to propagate an exception. The possible consequences depend on what action invoked the **Finalize** or **Adjust** operation:
  - 15 • For a **Finalize** invoked as part of an **assignment\_statement**, **Program\_Error** is raised at that point.
  - 16/2 • For an **Adjust** invoked as part of assignment operations other than those invoked as part of an **assignment\_statement**, other adjustments due to be performed might or might not be performed, and then **Program\_Error** is raised. During its propagation, finalization might or might not be applied to objects whose **Adjust** failed. For an **Adjust** invoked as part of an **assignment\_statement**, any other adjustments due to be performed are performed, and then **Program\_Error** is raised.
  - 17 • For a **Finalize** invoked as part of a call on an instance of **Unchecked\_Deallocation**, any other finalizations due to be performed are performed, and then **Program\_Error** is raised.
  - 17.1/1 • For a **Finalize** invoked as part of the finalization of the anonymous object created by a function call or **aggregate**, any other finalizations due to be performed are performed, and then **Program\_Error** is raised.
  - 17.2/1 • For a **Finalize** invoked due to reaching the end of the execution of a master, any other finalizations associated with the master are performed, and **Program\_Error** is raised immediately after leaving the master.
  - 18/2 • For a **Finalize** invoked by the transfer of control of an **exit\_statement**, **return statement**, **goto\_statement**, or **requeue\_statement**, **Program\_Error** is raised no earlier than after the finalization of the master being finalized when the exception occurred, and no later than the point where normal execution would have continued. Any other finalizations due to be performed up to that point are performed before raising **Program\_Error**.
  - 19 • For a **Finalize** invoked by a transfer of control that is due to raising an exception, any other finalizations due to be performed for the same master are performed; **Program\_Error** is raised immediately after leaving the master.
  - 20 • For a **Finalize** invoked by a transfer of control due to an abort or selection of a terminate alternative, the exception is ignored; any other finalizations due to be performed are performed.

#### NOTES

- 21 17 The rules of Section 10 imply that immediately prior to partition termination, **Finalize** operations are applied to library-level controlled objects (including those created by allocators of library-level access types, except those already finalized). This occurs after waiting for library-level tasks to terminate.

- |                                                                                                                                                                                                                                                                                                                                                                                                                                                             |    |
|-------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|----|
| 18 A constant is only constant between its initialization and finalization. Both initialization and finalization are allowed to change the value of a constant.                                                                                                                                                                                                                                                                                             | 22 |
| 19 Abort is deferred during certain operations related to controlled types, as explained in 9.8. Those rules prevent an abort from causing a controlled object to be left in an ill-defined state.                                                                                                                                                                                                                                                          | 23 |
| 20 The Finalize procedure is called upon finalization of a controlled object, even if Finalize was called earlier, either explicitly or as part of an assignment; hence, if a controlled type is visibly controlled (implying that its Finalize primitive is directly callable), or is nonlimited (implying that assignment is allowed), its Finalize procedure should be designed to have no ill effect if it is applied a second time to the same object. | 24 |

## Section 8: Visibility Rules

The rules defining the scope of declarations and the rules defining which identifiers, character\_literals, and operator\_symbols are visible at (or from) various places in the text of the program are described in this section. The formulation of these rules uses the notion of a declarative region.

As explained in Section 3, a declaration declares a view of an entity and associates a defining name with that view. The view comprises an identification of the viewed entity, and possibly additional properties. A usage name denotes a declaration. It also denotes the view declared by that declaration, and denotes the entity of that view. Thus, two different usage names might denote two different views of the same entity; in this case they denote the same entity.

### 8.1 Declarative Region

#### *Static Semantics*

For each of the following constructs, there is a portion of the program text called its *declarative region*, within which nested declarations can occur:

- any declaration, other than that of an enumeration type, that is not a completion of a previous declaration;
- a `block_statement`;
- a `loop_statement`;
- an `extended_return_statement`;
- an `accept_statement`;
- an `exception_handler`.

The declarative region includes the text of the construct together with additional text determined (recursively), as follows:

- If a declaration is included, so is its completion, if any.
- If the declaration of a library unit (including Standard — see 10.1.1) is included, so are the declarations of any child units (and their completions, by the previous rule). The child declarations occur after the declaration.
- If a `body_stub` is included, so is the corresponding `subunit`.
- If a `type_declaration` is included, then so is a corresponding `record_representation_clause`, if any.

The declarative region of a declaration is also called the *declarative region* of any view or entity declared by the declaration.

A declaration occurs *immediately within* a declarative region if this region is the innermost declarative region that encloses the declaration (the *immediately enclosing* declarative region), not counting the declarative region (if any) associated with the declaration itself.

A declaration is *local* to a declarative region if the declaration occurs immediately within the declarative region. An entity is *local* to a declarative region if the entity is declared by a declaration that is local to the declarative region.

- 15 A declaration is *global* to a declarative region if the declaration occurs immediately within another declarative region that encloses the declarative region. An entity is *global* to a declarative region if the entity is declared by a declaration that is global to the declarative region.

NOTES

16 1 The children of a parent library unit are inside the parent's declarative region, even though they do not occur inside the parent's declaration or body. This implies that one can use (for example) "P.Q" to refer to a child of P whose defining name is Q, and that after "use P;" Q can refer (directly) to that child.

17 2 As explained above and in 10.1.1, "Compilation Units - Library Units", all library units are descendants of Standard, and so are contained in the declarative region of Standard. They are *not* inside the declaration or body of Standard, but they *are* inside its declarative region.

18 3 For a declarative region that comes in multiple parts, the text of the declarative region does not contain any text that might appear between the parts. Thus, when a portion of a declarative region is said to extend from one place to another in the declarative region, the portion does not contain any text that might appear between the parts of the declarative region.

## 8.2 Scope of Declarations

- 1 For each declaration, the language rules define a certain portion of the program text called the *scope* of the declaration. The scope of a declaration is also called the scope of any view or entity declared by the declaration. Within the scope of an entity, and only there, there are places where it is legal to refer to the declared entity. These places are defined by the rules of visibility and overloading.

*Static Semantics*

- 2 The *immediate scope* of a declaration is a portion of the declarative region immediately enclosing the declaration. The immediate scope starts at the beginning of the declaration, except in the case of an overloadable declaration, in which case the immediate scope starts just after the place where the profile of the callable entity is determined (which is at the end of the *\_specification* for the callable entity, or at the end of the *generic\_instantiation* if an instance). The immediate scope extends to the end of the declarative region, with the following exceptions:

- 3 • The immediate scope of a *library\_item* includes only its semantic dependents.  
4 • The immediate scope of a declaration in the private part of a library unit does not include the visible part of any public descendant of that library unit.

- 5 The *visible part* of (a view of) an entity is a portion of the text of its declaration containing declarations that are visible from outside. The *private part* of (a view of) an entity that has a visible part contains all declarations within the declaration of (the view of) the entity, except those in the visible part; these are not visible from outside. Visible and private parts are defined only for these kinds of entities: callable entities, other program units, and composite types.

- 6 • The visible part of a view of a callable entity is its profile.  
7 • The visible part of a composite type other than a task or protected type consists of the declarations of all components declared (explicitly or implicitly) within the *type\_declaration*.  
8 • The visible part of a generic unit includes the *generic\_formal\_part*. For a generic package, it also includes the first list of *basic\_declarative\_items* of the *package\_specification*. For a generic subprogram, it also includes the profile.  
9 • The visible part of a package, task unit, or protected unit consists of declarations in the program unit's declaration other than those following the reserved word **private**, if any; see 7.1 and 12.7 for packages, 9.1 for task units, and 9.4 for protected units.

The scope of a declaration always contains the immediate scope of the declaration. In addition, for a given declaration that occurs immediately within the visible part of an outer declaration, or is a public child of an outer declaration, the scope of the given declaration extends to the end of the scope of the outer declaration, except that the scope of a `library_item` includes only its semantic dependents.

The scope of an `attribute_definition_clause` is identical to the scope of a declaration that would occur at the point of the `attribute_definition_clause`.

The immediate scope of a declaration is also the immediate scope of the entity or view declared by the declaration. Similarly, the scope of a declaration is also the scope of the entity or view declared by the declaration.

#### NOTES

4 There are notations for denoting visible declarations that are not directly visible. For example, `parameter_specifications` are in the visible part of a `subprogram_declaration` so that they can be used in named-notation calls appearing outside the called subprogram. For another example, declarations of the visible part of a package can be denoted by expanded names appearing outside the package, and can be made directly visible by a `use_clause`.

## 8.3 Visibility

The *visibility rules*, given below, determine which declarations are visible and directly visible at each place within a program. The visibility rules apply to both explicit and implicit declarations.

#### *Static Semantics*

A declaration is defined to be *directly visible* at places where a `name` consisting of only an identifier or `operator_symbol` is sufficient to denote the declaration; that is, no `selected_component` notation or special context (such as preceding `=>` in a named association) is necessary to denote the declaration. A declaration is defined to be *visible* wherever it is directly visible, as well as at other places where some `name` (such as a `selected_component`) can denote the declaration.

The syntactic category `direct_name` is used to indicate contexts where direct visibility is required. The syntactic category `selector_name` is used to indicate contexts where visibility, but not direct visibility, is required.

There are two kinds of direct visibility: *immediate visibility* and *use-visibility*. A declaration is immediately visible at a place if it is directly visible because the place is within its immediate scope. A declaration is use-visible if it is directly visible because of a `use_clause` (see 8.4). Both conditions can apply.

A declaration can be *hidden*, either from direct visibility, or from all visibility, within certain parts of its scope. Where *hidden from all visibility*, it is not visible at all (neither using a `direct_name` nor a `selector_name`). Where *hidden from direct visibility*, only direct visibility is lost; visibility using a `selector_name` is still possible.

Two or more declarations are *overloaded* if they all have the same defining name and there is a place where they are all directly visible.

The declarations of callable entities (including enumeration literals) are *overloadable*, meaning that overloading is allowed for them.

Two declarations are *homographs* if they have the same defining name, and, if both are overloadable, their profiles are type conformant. An inner declaration hides any outer homograph from direct visibility.

- 9/1 Two homographs are not generally allowed immediately within the same declarative region unless one *overrides* the other (see Legality Rules below). The only declarations that are *overridable* are the implicit declarations for predefined operators and inherited primitive subprograms. A declaration overrides another homograph that occurs immediately within the same declarative region in the following cases:
- 10/1 • A declaration that is not overridable overrides one that is overridable, regardless of which declaration occurs first;
  - 11 • The implicit declaration of an inherited operator overrides that of a predefined operator;
  - 12 • An implicit declaration of an inherited subprogram overrides a previous implicit declaration of an inherited subprogram.
  - 12.1/2 • If two or more homographs are implicitly declared at the same place:
    - 12.2/2 • If at least one is a subprogram that is neither a null procedure nor an abstract subprogram, and does not require overriding (see 3.9.3), then they override those that are null procedures, abstract subprograms, or require overriding. If more than one such homograph remains that is not thus overridden, then they are all hidden from all visibility.
    - 12.3/2 • Otherwise (all are null procedures, abstract subprograms, or require overriding), then any null procedure overrides all abstract subprograms and all subprograms that require overriding; if more than one such homograph remains that is not thus overridden, then if they are all fully conformant with one another, one is chosen arbitrarily; if not, they are all hidden from all visibility.
  - 13 • For an implicit declaration of a primitive subprogram in a generic unit, there is a copy of this declaration in an instance. However, a whole new set of primitive subprograms is implicitly declared for each type declared within the visible part of the instance. These new declarations occur immediately after the type declaration, and override the copied ones. The copied ones can be called only from within the instance; the new ones can be called only from outside the instance, although for tagged types, the body of a new one can be executed by a call to an old one.
- 14 A declaration is visible within its scope, except where hidden from all visibility, as follows:
- 15 • An overridden declaration is hidden from all visibility within the scope of the overriding declaration.
  - 16 • A declaration is hidden from all visibility until the end of the declaration, except:
    - 17 • For a record type or record extension, the declaration is hidden from all visibility only until the reserved word **record**;
    - 18/2 • For a **package\_declaration**, **generic\_package\_declaration**, or **subprogram\_body**, the declaration is hidden from all visibility only until the reserved word **is** of the declaration;
    - 18.1/2 • For a task declaration or protected declaration, the declaration is hidden from all visibility only until the reserved word **with** of the declaration if there is one, or the reserved word **is** of the declaration if there is no **with**.
  - 19 • If the completion of a declaration is a declaration, then within the scope of the completion, the first declaration is hidden from all visibility. Similarly, a **discriminant\_specification** or **parameter\_specification** is hidden within the scope of a corresponding **discriminant\_specification** or **parameter\_specification** of a corresponding completion, or of a corresponding **accept\_statement**.
  - 20/2 • The declaration of a library unit (including a **library\_unit\_renaming\_declaration**) is hidden from all visibility at places outside its declarative region that are not within the scope of a **nonlimited\_with\_clause** that mentions it. The limited view of a library package is hidden from all visibility at places that are not within the scope of a **limited\_with\_clause** that mentions it; in

addition, the limited view is hidden from all visibility within the declarative region of the package, as well as within the scope of any `nonlimited_with_clause` that mentions the package. Where the declaration of the limited view of a package is visible, any name that denotes the package denotes the limited view, including those provided by a package renaming.

- For each declaration or renaming of a generic unit as a child of some parent generic package, there is a corresponding declaration nested immediately within each instance of the parent. Such a nested declaration is hidden from all visibility except at places that are within the scope of a `with_clause` that mentions the child. 20.1/2

A declaration with a `defining_identifier` or `defining_operator_symbol` is immediately visible (and hence directly visible) within its immediate scope except where hidden from direct visibility, as follows: 21

- A declaration is hidden from direct visibility within the immediate scope of a homograph of the declaration, if the homograph occurs within an inner declarative region; 22
- A declaration is also hidden from direct visibility where hidden from all visibility. 23

An `attribute_definition_clause` is *visible* everywhere within its scope. 23.1/2

#### *Name Resolution Rules*

A `direct_name` shall resolve to denote a directly visible declaration whose defining name is the same as the `direct_name`. A `selector_name` shall resolve to denote a visible declaration whose defining name is the same as the `selector_name`. 24

These rules on visibility and direct visibility do not apply in a `context_clause`, a `parent_unit_name`, or a pragma that appears at the place of a `compilation_unit`. For those contexts, see the rules in 10.1.6, “Environment-Level Visibility Rules”. 25

#### *Legality Rules*

A non-overridable declaration is illegal if there is a homograph occurring immediately within the same declarative region that is visible at the place of the declaration, and is not hidden from all visibility by the non-overridable declaration. In addition, a type extension is illegal if somewhere within its immediate scope it has two visible components with the same name. Similarly, the `context_clause` for a compilation unit is illegal if it mentions (in a `with_clause`) some library unit, and there is a homograph of the library unit that is visible at the place of the compilation unit, and the homograph and the mentioned library unit are both declared immediately within the same declarative region. These rules also apply to dispatching operations declared in the visible part of an instance of a generic unit. However, they do not apply to other overloadable declarations in an instance; such declarations may have type conformant profiles in the instance, so long as the corresponding declarations in the generic were not type conformant. 26/2

#### NOTES

5 Visibility for compilation units follows from the definition of the environment in 10.1.4, except that it is necessary to apply a `with_clause` to obtain visibility to a `library_unit_declaration` or `library_unit_renaming_declaration`. 27

6 In addition to the visibility rules given above, the meaning of the occurrence of a `direct_name` or `selector_name` at a given place in the text can depend on the overloading rules (see 8.6). 28

7 Not all contexts where an `identifier`, `character_literal`, or `operator_symbol` are allowed require visibility of a corresponding declaration. Contexts where visibility is not required are identified by using one of these three syntactic categories directly in a syntax rule, rather than using `direct_name` or `selector_name`. 29

### 8.3.1 Overriding Indicators

- 1/2 An **overriding\_indicator** is used to declare that an operation is intended to override (or not override) an inherited operation.

#### *Syntax*

2/2      overriding\_indicator ::= [not] **overriding**

#### *Legality Rules*

- 3/2 If an **abstract\_subprogram\_declaration**, **null\_procedure\_declaration**, **subprogram\_body**, **subprogram\_body\_stub**, **subprogram\_renaming\_declaration**, **generic\_instantiation** of a subprogram, or **subprogram\_declaration** other than a protected subprogram has an **overriding\_indicator**, then:

- 4/2 • the operation shall be a primitive operation for some type;
- 5/2 • if the **overriding\_indicator** is **overriding**, then the operation shall override a homograph at the place of the declaration or body;
- 6/2 • if the **overriding\_indicator** is **not overriding**, then the operation shall not override any homograph (at any place).

- 7/2 In addition to the places where Legality Rules normally apply, these rules also apply in the private part of an instance of a generic unit.

#### NOTES

- 8/2 8 Rules for **overriding\_indicators** of task and protected entries and of protected subprograms are found in 9.5.2 and 9.4, respectively.

#### *Examples*

- 9/2 The use of **overriding\_indicators** allows the detection of errors at compile-time that otherwise might not be detected at all. For instance, we might declare a security queue derived from the Queue interface of 3.9.4 as:

```
10/2 type Security_Queue is new Queue with record ...;
11/2 overriding
12/2 procedure Append(Q : in out Security_Queue; Person : in Person_Name);
13/2 overriding
14/2 procedure Remove_First(Q : in out Security_Queue; Person : in Person_Name);
15/2 overriding
16/2 function Cur_Count(Q : in Security_Queue) return Natural;
17/2 overriding
18/2 function Max_Count(Q : in Security_Queue) return Natural;
19/2 not overriding
20/2 procedure Arrest(Q : in out Security_Queue; Person : in Person_Name);
```

- 16/2 The first four subprogram declarations guarantee that these subprograms will override the four subprograms inherited from the Queue interface. A misspelling in one of these subprograms will be detected by the implementation. Conversely, the declaration of Arrest guarantees that this is a new operation.

## 8.4 Use Clauses

A `use_package_clause` achieves direct visibility of declarations that appear in the visible part of a package; a `use_type_clause` achieves direct visibility of the primitive operators of a type.

### Syntax

```
use_clause ::= use_package_clause | use_type_clause
use_package_clause ::= use package_name {, package_name};
use_type_clause ::= use type subtype_mark {, subtype_mark};
```

### Legality Rules

A `package_name` of a `use_package_clause` shall denote a nonlimited view of a package.

### Static Semantics

For each `use_clause`, there is a certain region of text called the *scope* of the `use_clause`. For a `use_clause` within a `context_clause` of a `library_unit_declaration` or `library_unit_renaming_declaration`, the scope is the entire declarative region of the declaration. For a `use_clause` within a `context_clause` of a body, the scope is the entire body and any subunits (including multiply nested subunits). The scope does not include `context_clauses` themselves.

For a `use_clause` immediately within a declarative region, the scope is the portion of the declarative region starting just after the `use_clause` and extending to the end of the declarative region. However, the scope of a `use_clause` in the private part of a library unit does not include the visible part of any public descendant of that library unit.

A package is *named* in a `use_package_clause` if it is denoted by a `package_name` of that clause. A type is *named* in a `use_type_clause` if it is determined by a `subtype_mark` of that clause.

For each package named in a `use_package_clause` whose scope encloses a place, each declaration that occurs immediately within the declarative region of the package is *potentially use-visible* at this place if the declaration is visible at this place. For each type *T* or *TClass* named in a `use_type_clause` whose scope encloses a place, the declaration of each primitive operator of type *T* is potentially use-visible at this place if its declaration is visible at this place.

A declaration is *use-visible* if it is potentially use-visible, except in these naming-conflict cases:

- A potentially use-visible declaration is not use-visible if the place considered is within the immediate scope of a homograph of the declaration.
- Potentially use-visible declarations that have the same identifier are not use-visible unless each of them is an overloadable declaration.

### Dynamic Semantics

The elaboration of a `use_clause` has no effect.

### Examples

*Example of a use clause in a context clause:*

```
with Ada.Calendar; use Ada;
```

15    Example of a use type clause:

16    `use type Rational_Numbers.Rational; -- see 7.1  
Two_Thirds: Rational_Numbers.Rational := 2/3;`

## 8.5 Renaming Declarations

- 1    A renaming\_declaration declares another name for an entity, such as an object, exception, package, subprogram, entry, or generic unit. Alternatively, a subprogram\_renaming\_declaration can be the completion of a previous subprogram\_declaration.

*Syntax*

2    `renaming_declaration ::=  
    object_renaming_declaration  
  | exception_renaming_declaration  
  | package_renaming_declaration  
  | subprogram_renaming_declaration  
  | generic_renaming_declaration`

*Dynamic Semantics*

- 3    The elaboration of a renaming\_declaration evaluates the name that follows the reserved word renames and thereby determines the view and entity denoted by this name (the *renamed view* and *renamed entity*). A name that denotes the renaming\_declaration denotes (a new view of) the renamed entity.

**NOTES**

- 4    9 Renaming may be used to resolve name conflicts and to act as a shorthand. Renaming with a different identifier or operator\_symbol does not hide the old name; the new name and the old name need not be visible at the same places.
- 5    10 A task or protected object that is declared by an explicit object\_declaration can be renamed as an object. However, a single task or protected object cannot be renamed since the corresponding type is anonymous (meaning it has no nameable subtypes). For similar reasons, an object of an anonymous array or access type cannot be renamed.
- 6    11 A subtype defined without any additional constraint can be used to achieve the effect of renaming another subtype (including a task or protected subtype) as in

7    `subtype Mode is Ada.Text_IO.File_Mode;`

### 8.5.1 Object Renaming Declarations

- 1    An object\_renaming\_declaration is used to rename an object.

*Syntax*

2/2    `object_renaming_declaration ::=  
    defining_identifier : [null_exclusion] subtype_mark renames object_name;  
  | defining_identifier : access_definition renames object_name;`

*Name Resolution Rules*

- 3/2    The type of the object\_name shall resolve to the type determined by the subtype\_mark, or in the case where the type is defined by an access\_definition, to an anonymous access type. If the anonymous access type is an access-to-object type, the type of the object\_name shall have the same designated type as that of the access\_definition. If the anonymous access type is an access-to-subprogram type, the type of the object\_name shall have a designated profile that is type conformant with that of the access\_definition.

*Legality Rules*

- 4    The renamed entity shall be an object.

In the case where the type is defined by an `access_definition`, the type of the renamed object and the type defined by the `access_definition`: 4.1/2

- shall both be access-to-object types with statically matching designated subtypes and with both or neither being access-to-constant types; or 4.2/2
- shall both be access-to-subprogram types with subtype conformant designated profiles. 4.3/2

For an `object_renaming_declaration` with a `null_exclusion` or an `access_definition` that has a `null_exclusion`: 4.4/2

- if the `object_name` denotes a generic formal object of a generic unit *G*, and the `object_renaming_declaration` occurs within the body of *G* or within the body of a generic unit declared within the declarative region of *G*, then the declaration of the formal object of *G* shall have a `null_exclusion`; 4.5/2
- otherwise, the subtype of the `object_name` shall exclude null. In addition to the places where Legality Rules normally apply (see 12.3), this rule applies also in the private part of an instance of a generic unit. 4.6/2

The renamed entity shall not be a subcomponent that depends on discriminants of a variable whose nominal subtype is unconstrained, unless this subtype is indefinite, or the variable is constrained by its initial value. A slice of an array shall not be renamed if this restriction disallows renaming of the array. In addition to the places where Legality Rules normally apply, these rules apply also in the private part of an instance of a generic unit. These rules also apply for a renaming that appears in the body of a generic unit, with the additional requirement that even if the nominal subtype of the variable is indefinite, its type shall not be a descendant of an untagged generic formal derived type. 5/2

#### *Static Semantics*

An `object_renaming_declaration` declares a new view of the renamed object whose properties are identical to those of the renamed view. Thus, the properties of the renamed object are not affected by the `renaming_declaration`. In particular, its value and whether or not it is a constant are unaffected; similarly, the null exclusion or constraints that apply to an object are not affected by renaming (any constraint implied by the `subtype_mark` or `access_definition` of the `object_renaming_declaration` is ignored). 6/2

#### *Examples*

*Example of renaming an object:* 7

```
declare
 L : Person renames Leftmost_Person; -- see 3.10.1
begin
 L.Age := L.Age + 1;
end;
```

## 8.5.2 Exception Renaming Declarations

An `exception_renaming_declaration` is used to rename an exception. 1

#### *Syntax*

```
exception_renaming_declaration ::= defining_identifier : exception renames exception_name; 2
```

#### *Legality Rules*

The renamed entity shall be an exception. 3

*Static Semantics*

- 4 An exception\_renaming\_declaration declares a new view of the renamed exception.

*Examples*

- 5 Example of renaming an exception:

```
6 EOF : exception renames Ada.IO_Exceptions.End_Error; -- see A.13
```

### 8.5.3 Package Renaming Declarations

- 1 A package\_renaming\_declaration is used to rename a package.

*Syntax*

```
2 package_renaming_declaration ::=
3 package defining_program_unit_name renames package_name;
```

*Legality Rules*

- 3 The renamed entity shall be a package.

- 3.1/2 If the package\_name of a package\_renaming\_declaration denotes a limited view of a package *P*, then a name that denotes the package\_renaming\_declaration shall occur only within the immediate scope of the renaming or the scope of a with\_clause that mentions the package *P* or, if *P* is a nested package, the innermost library package enclosing *P*.

*Static Semantics*

- 4 A package\_renaming\_declaration declares a new view of the renamed package.

- 4.1/2 At places where the declaration of the limited view of the renamed package is visible, a name that denotes the package\_renaming\_declaration denotes a limited view of the package (see 10.1.1).

*Examples*

- 5 Example of renaming a package:

```
6 package TM renames Table_Manager;
```

### 8.5.4 Subprogram Renaming Declarations

- 1 A subprogram\_renaming\_declaration can serve as the completion of a subprogram\_declaration; such a renaming\_declaration is called a *renaming-as-body*. A subprogram\_renaming\_declaration that is not a completion is called a *renaming-as-declaration*, and is used to rename a subprogram (possibly an enumeration literal) or an entry.

*Syntax*

```
2/2 subprogram_renaming_declaration ::=
3 [overriding_indicator]
4 subprogram_specification renames callable_entity_name;
```

*Name Resolution Rules*

- 3 The expected profile for the callable\_entity\_name is the profile given in the subprogram\_specification.

*Legality Rules*

The profile of a renaming-as-declaration shall be mode-conformant with that of the renamed callable entity. 4

For a parameter or result subtype of the `subprogram_specification` that has an explicit `null_exclusion`: 4.1/2

- if the `callable_entity_name` denotes a generic formal subprogram of a generic unit *G*, and the `subprogram_renaming_declaration` occurs within the body of a generic unit *G* or within the body of a generic unit declared within the declarative region of the generic unit *G*, then the corresponding parameter or result subtype of the formal subprogram of *G* shall have a `null_exclusion`; 4.2/2
- otherwise, the subtype of the corresponding parameter or result type of the renamed callable entity shall exclude null. In addition to the places where Legality Rules normally apply (see 12.3), this rule applies also in the private part of an instance of a generic unit. 4.3/2

The profile of a renaming-as-body shall conform fully to that of the declaration it completes. If the renaming-as-body completes that declaration before the subprogram it declares is frozen, the profile shall be mode-conformant with that of the renamed callable entity and the subprogram it declares takes its convention from the renamed subprogram; otherwise, the profile shall be subtype-conformant with that of the renamed callable entity and the convention of the renamed subprogram shall not be Intrinsic. A renaming-as-body is illegal if the declaration occurs before the subprogram whose declaration it completes is frozen, and the renaming renames the subprogram itself, through one or more subprogram renaming declarations, none of whose subprograms has been frozen. 5/1

The `callable_entity_name` of a renaming shall not denote a subprogram that requires overriding (see 3.9.3). 5.1/2

The `callable_entity_name` of a renaming-as-body shall not denote an abstract subprogram. 5.2/2

A name that denotes a formal parameter of the `subprogram_specification` is not allowed within the `callable_entity_name`. 6

*Static Semantics*

A renaming-as-declaration declares a new view of the renamed entity. The profile of this new view takes its subtypes, parameter modes, and calling convention from the original profile of the callable entity, while taking the formal parameter names and default\_expressions from the profile given in the `subprogram_renaming_declaration`. The new view is a function or procedure, never an entry. 7

*Dynamic Semantics*

For a call to a subprogram whose body is given as a renaming-as-body, the execution of the renaming-as-body is equivalent to the execution of a `subprogram_body` that simply calls the renamed subprogram with its formal parameters as the actual parameters and, if it is a function, returns the value of the call. 7.1/1

For a call on a renaming of a dispatching subprogram that is overridden, if the overriding occurred before the renaming, then the body executed is that of the overriding declaration, even if the overriding declaration is not visible at the place of the renaming; otherwise, the inherited or predefined subprogram is called. 8

*Bounded (Run-Time) Errors*

If a subprogram directly or indirectly renames itself, then it is a bounded error to call that subprogram. Possible consequences are that `Program_Error` or `Storage_Error` is raised, or that the call results in infinite recursion. 8.1/1

## NOTES

- 9    12 A procedure can only be renamed as a procedure. A function whose `defining_designator` is either an identifier or an `operator_symbol` can be renamed with either an identifier or an `operator_symbol`; for renaming as an operator, the subprogram specification given in the `renaming_declaration` is subject to the rules given in 6.6 for operator declarations. Enumeration literals can be renamed as functions; similarly, `attribute_references` that denote functions (such as references to `Succ` and `Pred`) can be renamed as functions. An entry can only be renamed as a procedure; the new name is only allowed to appear in contexts that allow a procedure name. An entry of a family can be renamed, but an entry family cannot be renamed as a whole.
- 10    13 The operators of the root numeric types cannot be renamed because the types in the profile are anonymous, so the corresponding specifications cannot be written; the same holds for certain attributes, such as `Pos`.
- 11    14 Calls with the new name of a renamed entry are `procedure_call_statements` and are not allowed at places where the syntax requires an `entry_call_statement` in `conditional_and_timed_entry_calls`, nor in an `asynchronous_select`; similarly, the `Count` attribute is not available for the new name.
- 12    15 The primitiveness of a renaming-as-declaration is determined by its profile, and by where it occurs, as for any declaration of (a view of) a subprogram; primitiveness is not determined by the renamed view. In order to perform a dispatching call, the subprogram name has to denote a primitive subprogram, not a non-primitive renaming of a primitive subprogram.

*Examples**Examples of subprogram renaming declarations:*

```
procedure My_Write(C : in Character) renames Pool(K).Write; -- see 4.1.3
function Real_Plus(Left, Right : Real) return Real renames "+";
function Int_Plus(Left, Right : Integer) return Integer renames "+";
function Rouge return Color renames Red; -- see 3.5.1
function Rot return Color renames Red;
function Rosso return Color renames Rouge;
function Next(X : Color) return Color renames Color'Succ; -- see 3.5.1
```

*Example of a subprogram renaming declaration with new parameter names:*

```
function "*" (X, Y : Vector) return Real renames Dot_Product; -- see 6.1
```

*Example of a subprogram renaming declaration with a new default expression:*

```
function Minimum(L : Link := Head) return Cell renames Min_Cell; -- see 6.1
```

## 8.5.5 Generic Renaming Declarations

1    A `generic_renaming_declaration` is used to rename a generic unit.

*Syntax*

```
generic_renaming_declaration ::=
 generic package defining_program_unit_name renames generic_package_name;
 | generic procedure defining_program_unit_name renames generic_procedure_name;
 | generic function defining_program_unit_name renames generic_function_name;
```

*Legality Rules*

3    The renamed entity shall be a generic unit of the corresponding kind.

*Static Semantics*

4    A `generic_renaming_declaration` declares a new view of the renamed generic unit.

## NOTES

16 Although the properties of the new view are the same as those of the renamed view, the place where the generic\_renaming\_declaration occurs may affect the legality of subsequent renamings and instantiations that denote the generic\_renaming\_declaration, in particular if the renamed generic unit is a library unit (see 10.1.1).

*Examples**Example of renaming a generic unit:*

```
generic package Enum_IO renames Ada.Text_IO.Enumeration_IO; -- see A.10.10
```

## 8.6 The Context of Overload Resolution

Because declarations can be overloaded, it is possible for an occurrence of a usage name to have more than one possible interpretation; in most cases, ambiguity is disallowed. This clause describes how the possible interpretations resolve to the actual interpretation.

Certain rules of the language (the Name Resolution Rules) are considered “overloading rules”. If a possible interpretation violates an overloading rule, it is assumed not to be the intended interpretation; some other possible interpretation is assumed to be the actual interpretation. On the other hand, violations of non-overloading rules do not affect which interpretation is chosen; instead, they cause the construct to be illegal. To be legal, there usually has to be exactly one acceptable interpretation of a construct that is a “complete context”, not counting any nested complete contexts.

The syntax rules of the language and the visibility rules given in 8.3 determine the possible interpretations. Most type checking rules (rules that require a particular type, or a particular class of types, for example) are overloading rules. Various rules for the matching of formal and actual parameters are overloading rules.

*Name Resolution Rules*

Overload resolution is applied separately to each *complete context*, not counting inner complete contexts. Each of the following constructs is a *complete context*:

- A context\_item.
- A declarative\_item or declaration.
- A statement.
- A pragma\_argument\_association.
- The expression of a case\_statement.

An (overall) *interpretation* of a complete context embodies its meaning, and includes the following information about the constituents of the complete context, not including constituents of inner complete contexts:

- for each constituent of the complete context, to which syntactic categories it belongs, and by which syntax rules; and
- for each usage name, which declaration it denotes (and, therefore, which view and which entity it denotes); and
- for a complete context that is a declarative\_item, whether or not it is a completion of a declaration, and (if so) which declaration it completes.

A *possible interpretation* is one that obeys the syntax rules and the visibility rules. An *acceptable interpretation* is a possible interpretation that obeys the *overloading rules*, that is, those rules that specify an expected type or expected profile, or specify how a construct shall *resolve* or be *interpreted*.

- 15 The *interpretation* of a constituent of a complete context is determined from the overall interpretation of the complete context as a whole. Thus, for example, “interpreted as a `function_call`,” means that the construct’s interpretation says that it belongs to the syntactic category `function_call`.
- 16 Each occurrence of a usage name *denotes* the declaration determined by its interpretation. It also denotes the view declared by its denoted declaration, except in the following cases:
- 17/2 • If a usage name appears within the declarative region of a `type_declaration` and denotes that same `type_declaration`, then it denotes the *current instance* of the type (rather than the type itself); the current instance of a type is the object or value of the type that is associated with the execution that evaluates the usage name. This rule does not apply if the usage name appears within the `subtype_mark` of an `access_definition` for an access-to-object type, or within the subtype of a parameter or result of an access-to-subprogram type.
- 18 • If a usage name appears within the declarative region of a `generic_declaration` (but not within its `generic_formal_part`) and it denotes that same `generic_declaration`, then it denotes the *current instance* of the generic unit (rather than the generic unit itself). See also 12.3.
- 19 A usage name that denotes a view also denotes the entity of that view.
- 20/2 The *expected type* for a given `expression`, `name`, or other construct determines, according to the *type resolution rules* given below, the types considered for the construct during overload resolution. The type resolution rules provide support for class-wide programming, universal literals, dispatching operations, and anonymous access types:
- 21 • If a construct is expected to be of any type in a class of types, or of the universal or class-wide type for a class, then the type of the construct shall resolve to a type in that class or to a universal type that covers the class.
- 22 • If the expected type for a construct is a specific type  $T$ , then the type of the construct shall resolve either to  $T$ , or:
- 23 • to  $T\text{Class}$ ; or
  - 24 • to a universal type that covers  $T$ ; or
  - 25/2 • when  $T$  is a specific anonymous access-to-object type (see 3.10) with designated type  $D$ , to an access-to-object type whose designated type is  $D\text{Class}$  or is covered by  $D$ ; or
  - 25.1/2 • when  $T$  is an anonymous access-to-subprogram type (see 3.10), to an access-to-subprogram type whose designated profile is type-conformant with that of  $T$ .
- 26 In certain contexts, such as in a `subprogram_renaming_declaration`, the Name Resolution Rules define an *expected profile* for a given `name`; in such cases, the `name` shall resolve to the name of a callable entity whose profile is type conformant with the expected profile.

*Legality Rules*

- 27/2 When a construct is one that requires that its expected type be a *single* type in a given class, the type of the construct shall be determinable solely from the context in which the construct appears, excluding the construct itself, but using the requirement that it be in the given class. Furthermore, the context shall not be one that expects any type in some class that contains types of the given class; in particular, the construct shall not be the operand of a `type_conversion`.
- 28 A complete context shall have at least one acceptable interpretation; if there is exactly one, then that one is chosen.
- 29 There is a *preference* for the primitive operators (and `ranges`) of the root numeric types `root_integer` and `root_real`. In particular, if two acceptable interpretations of a constituent of a complete context differ only

in that one is for a primitive operator (or `range`) of the type `root_integer` or `root_real`, and the other is not, the interpretation using the primitive operator (or `range`) of the root numeric type is *preferred*.

For a complete context, if there is exactly one overall acceptable interpretation where each constituent's interpretation is the same as or preferred (in the above sense) over those in all other overall acceptable interpretations, then that one overall acceptable interpretation is chosen. Otherwise, the complete context is *ambiguous*. 30

A complete context other than a `pragma_argument_association` shall not be ambiguous. 31

A complete context that is a `pragma_argument_association` is allowed to be ambiguous (unless otherwise specified for the particular pragma), but only if every acceptable interpretation of the pragma argument is as a `name` that statically denotes a callable entity. Such a `name` denotes all of the declarations determined by its interpretations, and all of the views declared by these declarations. 32

#### NOTES

17 If a usage name has only one acceptable interpretation, then it denotes the corresponding entity. However, this does not mean that the usage name is necessarily legal since other requirements exist which are not considered for overload resolution; for example, the fact that an expression is static, whether an object is constant, mode and subtype conformance rules, freezing rules, order of elaboration, and so on. 33

Similarly, subtypes are not considered for overload resolution (the violation of a constraint does not make a program illegal but raises an exception during program execution). 34

## Section 9: Tasks and Synchronization

The execution of an Ada program consists of the execution of one or more *tasks*. Each task represents a separate thread of control that proceeds independently and concurrently between the points where it *interacts* with other tasks. The various forms of task interaction are described in this section, and include:

- the activation and termination of a task;
- a call on a protected subprogram of a *protected object*, providing exclusive read-write access, or concurrent read-only access to shared data;
- a call on an entry, either of another task, allowing for synchronous communication with that task, or of a protected object, allowing for asynchronous communication with one or more other tasks using that same protected object;
- a timed operation, including a simple delay statement, a timed entry call or accept, or a timed asynchronous select statement (see next item);
- an asynchronous transfer of control as part of an asynchronous select statement, where a task stops what it is doing and begins execution at a different point in response to the completion of an entry call or the expiration of a delay;
- an abort statement, allowing one task to cause the termination of another task.

In addition, tasks can communicate indirectly by reading and updating (unprotected) shared variables, presuming the access is properly synchronized through some other kind of task interaction.

### *Static Semantics*

The properties of a task are defined by a corresponding task declaration and `task_body`, which together define a program unit called a *task unit*.

### *Dynamic Semantics*

Over time, tasks proceed through various *states*. A task is initially *inactive*; upon activation, and prior to its *termination* it is either *blocked* (as part of some task interaction) or *ready* to run. While ready, a task competes for the available *execution resources* that it requires to run.

#### NOTES

1 Concurrent task execution may be implemented on multicomputers, multiprocessors, or with interleaved execution on a single physical processor. On the other hand, whenever an implementation can determine that the required semantic effects can be achieved when parts of the execution of a given task are performed by different physical processors acting in parallel, it may choose to perform them in this way.

## 9.1 Task Units and Task Objects

A task unit is declared by a *task declaration*, which has a corresponding `task_body`. A task declaration may be a `task_type_declaration`, in which case it declares a named task type; alternatively, it may be a `single_task_declaration`, in which case it defines an anonymous task type, as well as declaring a named task object of that type.

### *Syntax*

```
task_type_declaration ::=
 task type defining_identifier [known_discriminant_part] [is
 [new interface_list with]
 task_definition];
```

- 3/2        single\_task\_declaration ::=  
             task\_defining\_identifier [is  
             [new interface\_list with]  
             task\_definition];
- 4        task\_definition ::=  
             {task\_item}  
             [ private  
             {task\_item} ]  
             end [task\_identifier]
- 5/1        task\_item ::= entry\_declaration | aspect\_clause
- 6        task\_body ::=  
             task body defining\_identifier is  
             declarative\_part  
             begin  
             handled\_sequence\_of\_statements  
             end [task\_identifier];
- 7        If a task\_identifier appears at the end of a task\_definition or task\_body, it shall repeat the defining\_identifier.

*Legality Rules*

- 8/2        This paragraph was deleted.

*Static Semantics*

- 9        A task\_definition defines a task type and its first subtype. The first list of task\_items of a task\_definition, together with the known\_discriminant\_part, if any, is called the visible part of the task unit. The optional list of task\_items after the reserved word private is called the private part of the task unit.
- 9.1/1        For a task declaration without a task\_definition, a task\_definition without task\_items is assumed.
- 9.2/2        For a task declaration with an interface\_list, the task type inherits user-defined primitive subprograms from each progenitor type (see 3.9.4), in the same way that a derived type inherits user-defined primitive subprograms from its progenitor types (see 3.4). If the first parameter of a primitive inherited subprogram is of the task type or an access parameter designating the task type, and there is an entry\_declaration for a single entry with the same identifier within the task declaration, whose profile is type conformant with the prefixed view profile of the inherited subprogram, the inherited subprogram is said to be *implemented* by the conforming task entry.

*Legality Rules*

- 9.3/2        A task declaration requires a completion, which shall be a task\_body, and every task\_body shall be the completion of some task declaration.
- 9.4/2        Each interface\_subtype\_mark of an interface\_list appearing within a task declaration shall denote a limited interface type that is not a protected interface.
- 9.5/2        The prefixed view profile of an explicitly declared primitive subprogram of a tagged task type shall not be type conformant with any entry of the task type, if the first parameter of the subprogram is of the task type or is an access parameter designating the task type.
- 9.6/2        For each primitive subprogram inherited by the type declared by a task declaration, at most one of the following shall apply:

- the inherited subprogram is overridden with a primitive subprogram of the task type, in which case the overriding subprogram shall be subtype conformant with the inherited subprogram and not abstract; or 9.7/2
- the inherited subprogram is implemented by a single entry of the task type; in which case its prefixed view profile shall be subtype conformant with that of the task entry. 9.8/2

If neither applies, the inherited subprogram shall be a null procedure. In addition to the places where Legality Rules normally apply (see 12.3), these rules also apply in the private part of an instance of a generic unit. 9.9/2

#### *Dynamic Semantics*

The elaboration of a task declaration elaborates the `task_definition`. The elaboration of a `single_task_declaration` also creates an object of an (anonymous) task type. 10

The elaboration of a `task_definition` creates the task type and its first subtype; it also includes the elaboration of the `entry_declarations` in the given order. 11

As part of the initialization of a task object, any `aspect_clauses` and any per-object constraints associated with `entry_declarations` of the corresponding `task_definition` are elaborated in the given order. 12/1

The elaboration of a `task_body` has no effect other than to establish that tasks of the type can from then on be activated without failing the Elaboration\_Check. 13

The execution of a `task_body` is invoked by the activation of a task of the corresponding type (see 9.2). 14

The content of a task object of a given task type includes: 15

- The values of the discriminants of the task object, if any; 16
- An entry queue for each entry of the task object; 17
- A representation of the state of the associated task. 18

#### NOTES

2 Other than in an `access_definition`, the name of a task unit within the declaration or body of the task unit denotes the current instance of the unit (see 8.6), rather than the first subtype of the corresponding task type (and thus the name cannot be used as a `subtype_mark`). 19/2

3 The notation of a `selected_component` can be used to denote a discriminant of a task (see 4.1.3). Within a task unit, the name of a discriminant of the task type denotes the corresponding discriminant of the current instance of the unit. 20

4 A task type is a limited type (see 7.5), and hence precludes use of `assignment_statements` and predefined equality operators. If an application needs to store and exchange task identities, it can do so by defining an access type designating the corresponding task objects and by using access values for identification purposes. Assignment is available for such an access type as for any access type. Alternatively, if the implementation supports the Systems Programming Annex, the Identity attribute can be used for task identification (see C.7.1). 21/2

#### *Examples*

*Examples of declarations of task types:* 22

```
task type Server is
 entry Next_Work_Item(WI : in Work_Item);
 entry Shut_Down;
end Server;

task type Keyboard_Driver(ID : Keyboard_ID := New_ID) is
 new Serial_Device with -- see 3.9.4
 entry Read (C: out Character);
 entry Write(C : in Character);
end Keyboard_Driver;
```

25 Examples of declarations of single tasks:

```

26 task Controller is
27 entry Request(Level) (D : Item); -- a family of entries
28 end Controller;
29
27 task Parser is
28 entry Next_Lexeme(L : in Lexical_Element);
29 entry Next_Action(A : out Parser_Action);
30 end;
31
32 task User; -- has no entries

```

29 Examples of task objects:

```

30 Agent : Server;
31 Teletype : Keyboard_Driver(TTY_ID);
32 Pool : array(1 .. 10) of Keyboard_Driver;

```

31 Example of access type designating task objects:

```

32 type Keyboard is access Keyboard_Driver;
33 Terminal : Keyboard := new Keyboard_Driver(Term_ID);

```

## 9.2 Task Execution - Task Activation

### *Dynamic Semantics*

- 1 The execution of a task of a given task type consists of the execution of the corresponding `task_body`. The initial part of this execution is called the *activation* of the task; it consists of the elaboration of the `declarative_part` of the `task_body`. Should an exception be propagated by the elaboration of its `declarative_part`, the activation of the task is defined to have *failed*, and it becomes a completed task.
- 2 A task object (which represents one task) can be a part of a stand-alone object, of an object created by an allocator, or of an anonymous object of a limited type, or a coextension of one of these. All tasks that are part or coextensions of any of the stand-alone objects created by the elaboration of `object_declarations` (or `generic_associations` of formal objects of mode `in`) of a single declarative region are activated together. All tasks that are part or coextensions of a single object that is not a stand-alone object are activated together.
- 3 For the tasks of a given declarative region, the activations are initiated within the context of the `handled_sequence_of_statements` (and its associated `exception_handlers` if any — see 11.2), just prior to executing the statements of the `handled_sequence_of_statements`. For a package without an explicit body or an explicit `handled_sequence_of_statements`, an implicit body or an implicit `null_statement` is assumed, as defined in 7.2.
- 4 For tasks that are part or coextensions of a single object that is not a stand-alone object, activations are initiated after completing any initialization of the outermost object enclosing these tasks, prior to performing any other operation on the outermost object. In particular, for tasks that are part or coextensions of the object created by the evaluation of an allocator, the activations are initiated as the last step of evaluating the allocator, prior to returning the new access value. For tasks that are part or coextensions of an object that is the result of a function call, the activations are not initiated until after the function returns.
- 5 The task that created the new tasks and initiated their activations (the *activator*) is blocked until all of these activations complete (successfully or not). Once all of these activations are complete, if the activation of any of the tasks has failed (due to the propagation of an exception), `Tasking_Error` is raised in the activator, at the place at which it initiated the activations. Otherwise, the activator proceeds with its

execution normally. Any tasks that are aborted prior to completing their activation are ignored when determining whether to raise `Tasking_Error`.

Should the task that created the new tasks never reach the point where it would initiate the activations (due to an abort or the raising of an exception), the newly created tasks become terminated and are never activated.

## NOTES

5 An entry of a task can be called before the task has been activated.

6 If several tasks are activated together, the execution of any of these tasks need not await the end of the activation of the other tasks.

7 A task can become completed during its activation either because of an exception or because it is aborted (see 9.8).

6

7

8

9

*Examples*

*Example of task activation:*

```
procedure P is
 A, B : Server; -- elaborate the task objects A, B
 C : Server; -- elaborate the task object C
begin
 -- the tasks A, B, C are activated together before the first statement
 ...
end;
```

10

11

## 9.3 Task Dependence - Termination of Tasks

*Dynamic Semantics*

Each task (other than an environment task — see 10.2) *depends* on one or more masters (see 7.6.1), as follows:

- If the task is created by the evaluation of an `allocator` for a given access type, it depends on each master that includes the elaboration of the declaration of the ultimate ancestor of the given access type.
- If the task is created by the elaboration of an `object_declaration`, it depends on each master that includes this elaboration.
- Otherwise, the task depends on the master of the outermost object of which it is a part (as determined by the accessibility level of that object — see 3.10.2 and 7.6.1), as well as on any master whose execution includes that of the master of the outermost object.

Furthermore, if a task depends on a given master, it is defined to depend on the task that executes the master, and (recursively) on any master of that task.

A task is said to be *completed* when the execution of its corresponding `task_body` is completed. A task is said to be *terminated* when any finalization of the `task_body` has been performed (see 7.6.1). The first step of finalizing a master (including a `task_body`) is to wait for the termination of any tasks dependent on the master. The task executing the master is blocked until all the dependents have terminated. Any remaining finalization is then performed and the master is left.

Completion of a task (and the corresponding `task_body`) can occur when the task is blocked at a `select`-statement with an open `terminate_alternative` (see 9.7.1); the open `terminate_alternative` is selected if and only if the following conditions are satisfied:

- The task depends on some completed master; and

2

3

3.1/2

4

5

6/1

7/2

- 8     • Each task that depends on the master considered is either already terminated or similarly blocked at a `select_statement` with an open `terminate_alternative`.
- 9     When both conditions are satisfied, the task considered becomes completed, together with all tasks that depend on the master considered that are not yet completed.

## NOTES

10    8 The full view of a limited private type can be a task type, or can have subcomponents of a task type. Creation of an object of such a type creates dependences according to the full type.

11    9 An `object_renaming_declaration` defines a new view of an existing entity and hence creates no further dependence.

12    10 The rules given for the collective completion of a group of tasks all blocked on `select_statements` with open `terminate_alternatives` ensure that the collective completion can occur only when there are no remaining active tasks that could call one of the tasks being collectively completed.

13    11 If two or more tasks are blocked on `select_statements` with open `terminate_alternatives`, and become completed collectively, their finalization actions proceed concurrently.

14    12 The completion of a task can occur due to any of the following:

- the raising of an exception during the elaboration of the `declarative_part` of the corresponding `task_body`;
- the completion of the `handled_sequence_of_statements` of the corresponding `task_body`;
- the selection of an open `terminate_alternative` of a `select_statement` in the corresponding `task_body`;
- the abort of the task.

*Examples**Example of task dependence:*

```
20 declare
 type Global is access Server; -- see 9.1
 A, B : Server;
 G : Global;
begin
 -- activation of A and B
 declare
 type Local is access Server;
 X : Global := new Server; -- activation of X.all
 L : Local := new Server; -- activation of L.all
 C : Server;
 begin
 -- activation of C
 G := X; -- both G and X designate the same task object
 ...
 end; -- await termination of C and L.all (but not X.all)
 ...
end; -- await termination of A, B, and G.all
```

## 9.4 Protected Units and Protected Objects

A *protected object* provides coordinated access to shared data, through calls on its visible *protected operations*, which can be *protected subprograms* or *protected entries*. A *protected unit* is declared by a *protected declaration*, which has a corresponding *protected\_body*. A protected declaration may be a *protected\_type\_declaration*, in which case it declares a named protected type; alternatively, it may be a *single\_protected\_declaration*, in which case it defines an anonymous protected type, as well as declaring a named protected object of that type.

### Syntax

```

protected_type_declaration ::= 2/2
 protected type defining_identifier [known_discriminant_part] is
 [new interface_list with]
 protected_definition;

single_protected_declaration ::= 3/2
 protected defining_identifier is
 [new interface_list with]
 protected_definition;

protected_definition ::= 4
 { protected_operation_declaration }
 [private
 { protected_element_declaration }]
 end [protected_identifier]

protected_operation_declaration ::= subprogram_declaration 5/1
 | entry_declaration
 | aspect_clause

protected_element_declaration ::= protected_operation_declaration 6
 | component_declaration

protected_body ::= 7
 protected body defining_identifier is
 { protected_operation_item }
 end [protected_identifier];

protected_operation_item ::= subprogram_declaration 8/1
 | subprogram_body
 | entry_body
 | aspect_clause

```

If a *protected\_identifier* appears at the end of a *protected\_definition* or *protected\_body*, it shall repeat the *defining\_identifier*.

### Legality Rules

*This paragraph was deleted.*

1

2/2

3/2

4

5/1

6

7

8/1

9

10/2

11/2

### Static Semantics

A *protected\_definition* defines a protected type and its first subtype. The list of *protected\_operation\_declarations* of a *protected\_definition*, together with the *known\_discriminant\_part*, if any, is called the visible part of the protected unit. The optional list of *protected\_element\_declarations* after the reserved word **private** is called the private part of the protected unit.

- 11.1/2 For a protected declaration with an *interface\_list*, the protected type inherits user-defined primitive subprograms from each progenitor type (see 3.9.4), in the same way that a derived type inherits user-defined primitive subprograms from its progenitor types (see 3.4). If the first parameter of a primitive inherited subprogram is of the protected type or an access parameter designating the protected type, and there is a *protected\_operation\_declaration* for a protected subprogram or single entry with the same identifier within the protected declaration, whose profile is type conformant with the prefixed view profile of the inherited subprogram, the inherited subprogram is said to be *implemented* by the conforming protected subprogram or entry.

*Legality Rules*

- 11.2/2 A protected declaration requires a completion, which shall be a *protected\_body*, and every *protected\_body* shall be the completion of some protected declaration.
- 11.3/2 Each *interface\_subtype\_mark* of an *interface\_list* appearing within a protected declaration shall denote a limited interface type that is not a task interface.
- 11.4/2 The prefixed view profile of an explicitly declared primitive subprogram of a tagged protected type shall not be type conformant with any protected operation of the protected type, if the first parameter of the subprogram is of the protected type or is an access parameter designating the protected type.
- 11.5/2 For each primitive subprogram inherited by the type declared by a protected declaration, at most one of the following shall apply:
- 11.6/2 • the inherited subprogram is overridden with a primitive subprogram of the protected type, in which case the overriding subprogram shall be subtype conformant with the inherited subprogram and not abstract; or
  - 11.7/2 • the inherited subprogram is implemented by a protected subprogram or single entry of the protected type, in which case its prefixed view profile shall be subtype conformant with that of the protected subprogram or entry.
- 11.8/2 If neither applies, the inherited subprogram shall be a null procedure. In addition to the places where Legality Rules normally apply (see 12.3), these rules also apply in the private part of an instance of a generic unit.
- 11.9/2 If an inherited subprogram is implemented by a protected procedure or an entry, then the first parameter of the inherited subprogram shall be of mode **out** or **in out**, or an access-to-variable parameter.
- 11.10/2 If a protected subprogram declaration has an *overriding\_indicator*, then at the point of the declaration:
- 11.11/2 • if the *overriding\_indicator* is **overriding**, then the subprogram shall implement an inherited subprogram;
  - 11.12/2 • if the *overriding\_indicator* is **not overriding**, then the subprogram shall not implement any inherited subprogram.
- 11.13/2 In addition to the places where Legality Rules normally apply (see 12.3), these rules also apply in the private part of an instance of a generic unit.

*Dynamic Semantics*

- 12 The elaboration of a protected declaration elaborates the *protected\_definition*. The elaboration of a *single\_protected\_declaration* also creates an object of an (anonymous) protected type.
- 13 The elaboration of a *protected\_definition* creates the protected type and its first subtype; it also includes the elaboration of the *component\_declarations* and *protected\_operation\_declarations* in the given order.

|                                                                                                                                                                                                                                                                                                                                                                                |    |
|--------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|----|
| As part of the initialization of a protected object, any per-object constraints (see 3.8) are elaborated.                                                                                                                                                                                                                                                                      | 14 |
| The elaboration of a <b>protected_body</b> has no other effect than to establish that protected operations of the type can from then on be called without failing the Elaboration_Check.                                                                                                                                                                                       | 15 |
| The content of an object of a given protected type includes:                                                                                                                                                                                                                                                                                                                   | 16 |
| <ul style="list-style-type: none"> <li>• The values of the components of the protected object, including (implicitly) an entry queue for each entry declared for the protected object;</li> <li>• A representation of the state of the execution resource <i>associated</i> with the protected object (one such resource is associated with each protected object).</li> </ul> | 17 |
| The execution resource associated with a protected object has to be acquired to read or update any components of the protected object; it can be acquired (as part of a protected action — see 9.5.1) either for concurrent read-only access, or for exclusive read-write access.                                                                                              | 19 |
| As the first step of the <i>finalization</i> of a protected object, each call remaining on any entry queue of the object is removed from its queue and Program_Error is raised at the place of the corresponding <b>entry-call_statement</b> .                                                                                                                                 | 20 |

#### *Bounded (Run-Time) Errors*

It is a bounded error to call an entry or subprogram of a protected object after that object is finalized. If the error is detected, Program\_Error is raised. Otherwise, the call proceeds normally, which may leave a task queued forever.

#### NOTES

- 13 Within the declaration or body of a protected unit other than in an **access\_definition**, the name of the protected unit denotes the current instance of the unit (see 8.6), rather than the first subtype of the corresponding protected type (and thus the name cannot be used as a **subtype\_mark**).
- 14 A **selected\_component** can be used to denote a discriminant of a protected object (see 4.1.3). Within a protected unit, the name of a discriminant of the protected type denotes the corresponding discriminant of the current instance of the unit.
- 15 A protected type is a limited type (see 7.5), and hence precludes use of **assignment\_statements** and predefined equality operators.
- 16 The bodies of the protected operations given in the **protected\_body** define the actions that take place upon calls to the protected operations.
- 17 The declarations in the private part are only visible within the private part and the body of the protected unit.

#### *Examples*

*Example of declaration of protected type and corresponding body:*

```

protected type Resource is
 entry Seize;
 procedure Release;
private
 Busy : Boolean := False;
end Resource;

protected body Resource is
 entry Seize when not Busy is
 begin
 Busy := True;
 end Seize;
 procedure Release is
 begin
 Busy := False;
 end Release;
end Resource;

```

Example of a single protected declaration and corresponding body:

```

30 protected Shared_Array is
31 -- Index, Item, and Item_Array are global types
32 function Component (N : in Index) return Item;
33 procedure Set_Component(N : in Index; E : in Item);
34 private
35 Table : Item_Array(Index) := (others => Null_Item);
36 end Shared_Array;

37 protected body Shared_Array is
38 function Component(N : in Index) return Item is
39 begin
40 return Table(N);
41 end Component;
42
43 procedure Set_Component(N : in Index; E : in Item) is
44 begin
45 Table(N) := E;
46 end Set_Component;
47 end Shared_Array;
```

Examples of protected objects:

```

48 Control : Resource;
49 Flags : array(1 .. 100) of Resource;
```

## 9.5 Intertask Communication

- 1 The primary means for intertask communication is provided by calls on entries and protected subprograms. Calls on protected subprograms allow coordinated access to shared data objects. Entry calls allow for blocking the caller until a given condition is satisfied (namely, that the corresponding entry is open — see 9.5.3), and then communicating data or control information directly with another task or indirectly via a shared protected object.

### Static Semantics

- 2 Any call on an entry or on a protected subprogram identifies a *target object* for the operation, which is either a task (for an entry call) or a protected object (for an entry call or a protected subprogram call). The target object is considered an implicit parameter to the operation, and is determined by the operation name (or prefix) used in the call on the operation, as follows:
- 3 • If it is a *direct\_name* or expanded name that denotes the declaration (or body) of the operation, then the target object is implicitly specified to be the current instance of the task or protected unit immediately enclosing the operation; such a call is defined to be an *internal call*;
  - 4 • If it is a *selected\_component* that is not an expanded name, then the target object is explicitly specified to be the task or protected object denoted by the prefix of the name; such a call is defined to be an *external call*;
  - 5 • If the name or prefix is a dereference (implicit or explicit) of an access-to-protected-subprogram value, then the target object is determined by the prefix of the *Access attribute reference* that produced the access value originally, and the call is defined to be an *external call*;
  - 6 • If the name or prefix denotes a *subprogram\_renaming\_declaration*, then the target object is as determined by the name of the renamed entity.
- 7 A corresponding definition of target object applies to a *queue\_statement* (see 9.5.4), with a corresponding distinction between an *internal queue* and an *external queue*.

*Legality Rules*

The view of the target protected object associated with a call of a protected procedure or entry shall be a variable. 7.1/2

*Dynamic Semantics*

Within the body of a protected operation, the current instance (see 8.6) of the immediately enclosing protected unit is determined by the target object specified (implicitly or explicitly) in the call (or requeue) on the protected operation. 8

Any call on a protected procedure or entry of a target protected object is defined to be an update to the object, as is a requeue on such an entry. 9

### 9.5.1 Protected Subprograms and Protected Actions

A *protected subprogram* is a subprogram declared immediately within a *protected\_definition*. Protected procedures provide exclusive read-write access to the data of a protected object; protected functions provide concurrent read-only access to the data. 1

*Static Semantics*

Within the body of a protected function (or a function declared immediately within a *protected\_body*), the current instance of the enclosing protected unit is defined to be a constant (that is, its subcomponents may be read but not updated). Within the body of a protected procedure (or a procedure declared immediately within a *protected\_body*), and within an *entry\_body*, the current instance is defined to be a variable (updating is permitted). 2

*Dynamic Semantics*

For the execution of a call on a protected subprogram, the evaluation of the name or prefix and of the parameter associations, and any assigning back of **in out** or **out** parameters, proceeds as for a normal subprogram call (see 6.4). If the call is an internal call (see 9.5), the body of the subprogram is executed as for a normal subprogram call. If the call is an external call, then the body of the subprogram is executed as part of a new *protected action* on the target protected object; the protected action completes after the body of the subprogram is executed. A protected action can also be started by an entry call (see 9.5.3). 3

A new protected action is not started on a protected object while another protected action on the same protected object is underway, unless both actions are the result of a call on a protected function. This rule is expressible in terms of the execution resource associated with the protected object: 4

- *Starting* a protected action on a protected object corresponds to *acquiring* the execution resource associated with the protected object, either for concurrent read-only access if the protected action is for a call on a protected function, or for exclusive read-write access otherwise;
- *Completing* the protected action corresponds to *releasing* the associated execution resource. 6

After performing an operation on a protected object other than a call on a protected function, but prior to completing the associated protected action, the entry queues (if any) of the protected object are serviced (see 9.5.3). 7

*Bounded (Run-Time) Errors*

During a protected action, it is a bounded error to invoke an operation that is *potentially blocking*. The following are defined to be potentially blocking operations: 8

- a *select\_statement*; 9

- 10     • an accept\_statement;  
 11    • an entry\_call\_statement;  
 12    • a delay\_statement;  
 13    • an abort\_statement;  
 14    • task creation or activation;  
 15    • an external call on a protected subprogram (or an external requeue) with the same target object  
       as that of the protected action;  
 16    • a call on a subprogram whose body contains a potentially blocking operation.

17 If the bounded error is detected, Program\_Error is raised. If not detected, the bounded error might result in deadlock or a (nested) protected action on the same target object.

18 Certain language-defined subprograms are potentially blocking. In particular, the subprograms of the language-defined input-output packages that manipulate files (implicitly or explicitly) are potentially blocking. Other potentially blocking subprograms are identified where they are defined. When not specified as potentially blocking, a language-defined subprogram is nonblocking.

#### NOTES

19    18 If two tasks both try to start a protected action on a protected object, and at most one is calling a protected function, then only one of the tasks can proceed. Although the other task cannot proceed, it is not considered blocked, and it might be consuming processing resources while it awaits its turn. There is no language-defined ordering or queuing presumed for tasks competing to start a protected action — on a multiprocessor such tasks might use busy-waiting; for monoprocessor considerations, see D.3, “Priority Ceiling Locking”.

20    19 The body of a protected unit may contain declarations and bodies for local subprograms. These are not visible outside the protected unit.

21    20 The body of a protected function can contain internal calls on other protected functions, but not protected procedures, because the current instance is a constant. On the other hand, the body of a protected procedure can contain internal calls on both protected functions and procedures.

22    21 From within a protected action, an internal call on a protected subprogram, or an external call on a protected subprogram with a different target object is not considered a potentially blocking operation.

22.1/2   22 The pragma Detect\_Blocking may be used to ensure that all executions of potentially blocking operations during a protected action raise Program\_Error. See H.5.

#### *Examples*

23    *Examples of protected subprogram calls (see 9.4):*

```
24 Shared_Array.Set_Component (N, E);
 E := Shared_Array.Component (M);
 Control.Release;
```

## 9.5.2 Entries and Accept Statements

1    Entry\_declarations, with the corresponding entry\_bodies or accept\_statements, are used to define potentially queued operations on tasks and protected objects.

#### *Syntax*

```
2/2 entry_declaration ::=
 [overriding_indicator]
 entry defining_identifier [(discrete_subtype_definition)] parameter_profile;
```

|                                                                                                                                                                                                                                     |        |
|-------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|--------|
| accept_statement ::=                                                                                                                                                                                                                | 3      |
| accept <i>entry_direct_name</i> [( <i>entry_index</i> )] parameter_profile [do<br>handled_sequence_of_statements<br>end [ <i>entry_identifier</i> ]];                                                                               |        |
| <i>entry_index</i> ::= expression                                                                                                                                                                                                   | 4      |
| <i>entry_body</i> ::=                                                                                                                                                                                                               | 5      |
| <i>entry</i> defining_identifier <i>entry_body_formal_part</i> <i>entry_barrier</i> is<br>declarative_part                                                                                                                          |        |
| begin                                                                                                                                                                                                                               |        |
| handled_sequence_of_statements                                                                                                                                                                                                      |        |
| end [ <i>entry_identifier</i> ];                                                                                                                                                                                                    |        |
| <i>entry_body_formal_part</i> ::= [( <i>entry_index_specification</i> )] parameter_profile                                                                                                                                          | 6      |
| <i>entry_barrier</i> ::= when condition                                                                                                                                                                                             | 7      |
| <i>entry_index_specification</i> ::= for defining_identifier in discrete_subtype_definition                                                                                                                                         | 8      |
| If an <i>entry_identifier</i> appears at the end of an accept_statement, it shall repeat the <i>entry_direct_name</i> . If an <i>entry_identifier</i> appears at the end of an entry_body, it shall repeat the defining_identifier. | 9      |
| An entry_declaration is allowed only in a protected or task declaration.                                                                                                                                                            | 10     |
| An overriding_indicator is not allowed in an entry_declaration that includes a discrete_subtype_definition.                                                                                                                         | 10.1/2 |

*Name Resolution Rules*

In an accept\_statement, the expected profile for the *entry\_direct\_name* is that of the entry\_declaration; the expected type for an *entry\_index* is that of the subtype defined by the discrete\_subtype\_definition of the corresponding entry\_declaration.

Within the handled\_sequence\_of\_statements of an accept\_statement, if a selected\_component has a prefix that denotes the corresponding entry\_declaration, then the entity denoted by the prefix is the accept\_statement, and the selected\_component is interpreted as an expanded name (see 4.1.3); the selector\_name of the selected\_component has to be the identifier for some formal parameter of the accept\_statement.

*Legality Rules*

An entry\_declaration in a task declaration shall not contain a specification for an access parameter (see 3.10).

- If an entry\_declaration has an overriding\_indicator, then at the point of the declaration:
- if the overriding\_indicator is overriding, then the entry shall implement an inherited subprogram;
  - if the overriding\_indicator is not overriding, then the entry shall not implement any inherited subprogram.

In addition to the places where Legality Rules normally apply (see 12.3), these rules also apply in the private part of an instance of a generic unit.

For an accept\_statement, the innermost enclosing body shall be a task\_body, and the *entry\_direct\_name* shall denote an entry\_declaration in the corresponding task declaration; the profile of the accept\_statement shall conform fully to that of the corresponding entry\_declaration. An accept\_statement shall

have a parenthesized `entry_index` if and only if the corresponding `entry_declaration` has a `discrete_subtype_definition`.

- 15 An `accept_statement` shall not be within another `accept_statement` that corresponds to the same `entry_declaration`, nor within an `asynchronous_select` inner to the enclosing `task_body`.
- 16 An `entry_declaration` of a protected unit requires a completion, which shall be an `entry_body`, and every `entry_body` shall be the completion of an `entry_declaration` of a protected unit. The profile of the `entry_body` shall conform fully to that of the corresponding declaration.
- 17 An `entry_body_formal_part` shall have an `entry_index_specification` if and only if the corresponding `entry_declaration` has a `discrete_subtype_definition`. In this case, the `discrete_subtype_definitions` of the `entry_declaration` and the `entry_index_specification` shall fully conform to one another (see 6.3.1).
- 18 A name that denotes a formal parameter of an `entry_body` is not allowed within the `entry_barrier` of the `entry_body`.

#### *Static Semantics*

- 19 The parameter modes defined for parameters in the `parameter_profile` of an `entry_declaration` are the same as for a `subprogram_declaration` and have the same meaning (see 6.2).
- 20 An `entry_declaration` with a `discrete_subtype_definition` (see 3.6) declares a *family* of distinct entries having the same profile, with one such entry for each value of the `entry_index_subtype` defined by the `discrete_subtype_definition`. A name for an entry of a family takes the form of an `indexed_component`, where the prefix denotes the `entry_declaration` for the family, and the index value identifies the entry within the family. The term *single entry* is used to refer to any entry other than an entry of an entry family.
- 21 In the `entry_body` for an entry family, the `entry_index_specification` declares a named constant whose subtype is the `entry_index_subtype` defined by the corresponding `entry_declaration`; the value of the *named entry index* identifies which entry of the family was called.

#### *Dynamic Semantics*

- 22 The elaboration of an `entry_declaration` for an entry family consists of the elaboration of the `discrete_subtype_definition`, as described in 3.8. The elaboration of an `entry_declaration` for a single entry has no effect.
- 23 The actions to be performed when an entry is called are specified by the corresponding `accept_statements` (if any) for an entry of a task unit, and by the corresponding `entry_body` for an entry of a protected unit.
- 24 For the execution of an `accept_statement`, the `entry_index`, if any, is first evaluated and converted to the `entry_index_subtype`; this index value identifies which entry of the family is to be accepted. Further execution of the `accept_statement` is then blocked until a caller of the corresponding entry is selected (see 9.5.3), whereupon the `handled_sequence_of_statements`, if any, of the `accept_statement` is executed, with the formal parameters associated with the corresponding actual parameters of the selected entry call. Upon completion of the `handled_sequence_of_statements`, the `accept_statement` completes and is left. When an exception is propagated from the `handled_sequence_of_statements` of an `accept_statement`, the same exception is also raised by the execution of the corresponding `entry_call_statement`.
- 25 The above interaction between a calling task and an accepting task is called a *rendezvous*. After a rendezvous, the two tasks continue their execution independently.

An `entry_body` is executed when the condition of the `entry_barrier` evaluates to True and a caller of the corresponding single entry, or entry of the corresponding entry family, has been selected (see 9.5.3). For the execution of the `entry_body`, the `declarative_part` of the `entry_body` is elaborated, and the `handled_sequence_of_statements` of the body is executed, as for the execution of a `subprogram_body`. The value of the named entry index, if any, is determined by the value of the entry index specified in the `entry_name` of the selected entry call (or intermediate `requeue_statement` — see 9.5.4).

## NOTES

- 23 A task entry has corresponding `accept_statements` (zero or more), whereas a protected entry has a corresponding `entry_body` (exactly one). 27
- 24 A consequence of the rule regarding the allowed placements of `accept_statements` is that a task can execute `accept_statements` only for its own entries. 28
- 25 A return statement (see 6.5) or a `requeue_statement` (see 9.5.4) may be used to complete the execution of an `accept_statement` or an `entry_body`. 29/2
- 26 The condition in the `entry_barrier` may reference anything visible except the formal parameters of the entry. This includes the entry index (if any), the components (including discriminants) of the protected object, the Count attribute of an entry of that protected object, and data global to the protected unit. 30

The restriction against referencing the formal parameters within an `entry_barrier` ensures that all calls of the same entry see the same barrier value. If it is necessary to look at the parameters of an entry call before deciding whether to handle it, the `entry_barrier` can be “`when True`” and the caller can be requeued (on some private entry) when its parameters indicate that it cannot be handled immediately.

*Examples**Examples of entry declarations:*

```
entry Read(V : out Item);
entry Seize;
entry Request(Level) (D : Item); -- a family of entries
```

*Examples of accept statements:*

```
accept Shut_Down;
accept Read(V : out Item) do
 V := Local_Item;
end Read;
accept Request(Low) (D : Item) do
 ...
end Request;
```

### 9.5.3 Entry Calls

An `entry_call_statement` (an *entry call*) can appear in various contexts. A *simple* entry call is a stand-alone statement that represents an unconditional call on an entry of a target task or a protected object. Entry calls can also appear as part of `select_statements` (see 9.7).

*Syntax*

```
entry_call_statement ::= entry_name [actual_parameter_part];
```

*Name Resolution Rules*

The `entry_name` given in an `entry_call_statement` shall resolve to denote an entry. The rules for parameter associations are the same as for subprogram calls (see 6.4 and 6.4.1).

*Static Semantics*

- 4 The `entry_name` of an `entry_call_statement` specifies (explicitly or implicitly) the target object of the call, the entry or entry family, and the entry index, if any (see 9.5).

*Dynamic Semantics*

- 5 Under certain circumstances (detailed below), an entry of a task or protected object is checked to see whether it is *open* or *closed*:

- 6 • An entry of a task is open if the task is blocked on an `accept_statement` that corresponds to the entry (see 9.5.2), or on a `selective_accept` (see 9.7.1) with an open `accept_alternative` that corresponds to the entry; otherwise it is closed.
- 7 • An entry of a protected object is open if the condition of the `entry_barrier` of the corresponding `entry_body` evaluates to True; otherwise it is closed. If the evaluation of the condition propagates an exception, the exception `Program_Error` is propagated to all current callers of all entries of the protected object.
- 8 For the execution of an `entry_call_statement`, evaluation of the `name` and of the parameter associations is as for a subprogram call (see 6.4). The entry call is then *issued*: For a call on an entry of a protected object, a new protected action is started on the object (see 9.5.1). The named entry is checked to see if it is open; if open, the entry call is said to be *selected immediately*, and the execution of the call proceeds as follows:
  - 9 • For a call on an open entry of a task, the accepting task becomes ready and continues the execution of the corresponding `accept_statement` (see 9.5.2).
  - 10 • For a call on an open entry of a protected object, the corresponding `entry_body` is executed (see 9.5.2) as part of the protected action.
- 11 If the `accept_statement` or `entry_body` completes other than by a requeue (see 9.5.4), return is made to the caller (after servicing the entry queues — see below); any necessary assigning back of formal to actual parameters occurs, as for a subprogram call (see 6.4.1); such assignments take place outside of any protected action.
- 12 If the named entry is closed, the entry call is added to an *entry queue* (as part of the protected action, for a call on a protected entry), and the call remains queued until it is selected or cancelled; there is a separate (logical) entry queue for each entry of a given task or protected object (including each entry of an entry family).
- 13 When a queued call is *selected*, it is removed from its entry queue. Selecting a queued call from a particular entry queue is called *servicing* the entry queue. An entry with queued calls can be serviced under the following circumstances:
  - 14 • When the associated task reaches a corresponding `accept_statement`, or a `selective_accept` with a corresponding open `accept_alternative`;
  - 15 • If after performing, as part of a protected action on the associated protected object, an operation on the object other than a call on a protected function, the entry is checked and found to be open.
- 16 If there is at least one call on a queue corresponding to an open entry, then one such call is selected according to the *entry queuing policy* in effect (see below), and the corresponding `accept_statement` or `entry_body` is executed as above for an entry call that is selected immediately.
- 17 The entry queuing policy controls selection among queued calls both for task and protected entry queues. The default entry queuing policy is to select calls on a given entry queue in order of arrival. If calls from two or more queues are simultaneously eligible for selection, the default entry queuing policy does not specify which queue is serviced first. Other entry queuing policies can be specified by pragmas (see D.4).

For a protected object, the above servicing of entry queues continues until there are no open entries with queued calls, at which point the protected action completes.

For an entry call that is added to a queue, and that is not the triggering\_statement of an asynchronous\_select (see 9.7.4), the calling task is blocked until the call is cancelled, or the call is selected and a corresponding accept\_statement or entry\_body completes without requeueing. In addition, the calling task is blocked during a rendezvous.

An attempt can be made to cancel an entry call upon an abort (see 9.8) and as part of certain forms of select\_statement (see 9.7.2, 9.7.3, and 9.7.4). The cancellation does not take place until a point (if any) when the call is on some entry queue, and not protected from cancellation as part of a requeue (see 9.5.4); at such a point, the call is removed from the entry queue and the call completes due to the cancellation. The cancellation of a call on an entry of a protected object is a protected action, and as such cannot take place while any other protected action is occurring on the protected object. Like any protected action, it includes servicing of the entry queues (in case some entry barrier depends on a Count attribute).

A call on an entry of a task that has already completed its execution raises the exception Tasking\_Error at the point of the call; similarly, this exception is raised at the point of the call if the called task completes its execution or becomes abnormal before accepting the call or completing the rendezvous (see 9.8). This applies equally to a simple entry call and to an entry call as part of a select\_statement.

#### *Implementation Permissions*

An implementation may perform the sequence of steps of a protected action using any thread of control; it need not be that of the task that started the protected action. If an entry\_body completes without requeueing, then the corresponding calling task may be made ready without waiting for the entire protected action to complete.

When the entry of a protected object is checked to see whether it is open, the implementation need not reevaluate the condition of the corresponding entry\_barrier if no variable or attribute referenced by the condition (directly or indirectly) has been altered by the execution (or cancellation) of a protected procedure or entry call on the object since the condition was last evaluated.

An implementation may evaluate the conditions of all entry\_barriers of a given protected object any time any entry of the object is checked to see if it is open.

When an attempt is made to cancel an entry call, the implementation need not make the attempt using the thread of control of the task (or interrupt) that initiated the cancellation; in particular, it may use the thread of control of the caller itself to attempt the cancellation, even if this might allow the entry call to be selected in the interim.

#### NOTES

27 If an exception is raised during the execution of an entry\_body, it is propagated to the corresponding caller (see 11.4).

28 For a call on a protected entry, the entry is checked to see if it is open prior to queuing the call, and again thereafter if its Count attribute (see 9.9) is referenced in some entry barrier.

29 In addition to simple entry calls, the language permits timed, conditional, and asynchronous entry calls (see 9.7.2, 9.7.3, and see 9.7.4).

30 The condition of an entry\_barrier is allowed to be evaluated by an implementation more often than strictly necessary, even if the evaluation might have side effects. On the other hand, an implementation need not reevaluate the condition if nothing it references was updated by an intervening protected action on the protected object, even if the condition references some global variable that might have been updated by an action performed from outside of a protected action.

*Examples*30    *Examples of entry calls:*

|    |                                      |            |
|----|--------------------------------------|------------|
| 31 | Agent.Shut_Down;                     | -- see 9.1 |
|    | Parser.Next_Lexeme (E) ;             | -- see 9.1 |
|    | Pool(5).Read(Next_Char) ;            | -- see 9.1 |
|    | Controller.Request(Low)(Some_Item) ; | -- see 9.1 |
|    | Flags(3).Seize;                      | -- see 9.4 |

**9.5.4 Requeue Statements**1    A **requeue\_statement** can be used to complete an **accept\_statement** or **entry\_body**, while redirecting the corresponding entry call to a new (or the same) entry queue. Such a *requeue* can be performed with or without allowing an intermediate cancellation of the call, due to an abort or the expiration of a delay.*Syntax*2    **requeue\_statement ::= requeue entry\_name [with abort];***Name Resolution Rules*3    The **entry\_name** of a **requeue\_statement** shall resolve to denote an entry (the *target entry*) that either has no parameters, or that has a profile that is type conformant (see 6.3.1) with the profile of the innermost enclosing **entry\_body** or **accept\_statement**.*Legality Rules*

- 4    A **requeue\_statement** shall be within a callable construct that is either an **entry\_body** or an **accept\_statement**, and this construct shall be the innermost enclosing body or callable construct.
- 5    If the target entry has parameters, then its profile shall be subtype conformant with the profile of the innermost enclosing callable construct.
- 6    In a **requeue\_statement** of an **accept\_statement** of some task unit, either the target object shall be a part of a formal parameter of the **accept\_statement**, or the accessibility level of the target object shall not be equal to or statically deeper than any enclosing **accept\_statement** of the task unit. In a **requeue\_statement** of an **entry\_body** of some protected unit, either the target object shall be a part of a formal parameter of the **entry\_body**, or the accessibility level of the target object shall not be statically deeper than that of the **entry\_declaration**.

*Dynamic Semantics*

- 7    The execution of a **requeue\_statement** proceeds by first evaluating the **entry\_name**, including the prefix identifying the target task or protected object and the expression identifying the entry within an entry family, if any. The **entry\_body** or **accept\_statement** enclosing the **requeue\_statement** is then completed, finalized, and left (see 7.6.1).
- 8    For the execution of a **requeue** on an entry of a target task, after leaving the enclosing callable construct, the named entry is checked to see if it is open and the queued call is either selected immediately or queued, as for a normal entry call (see 9.5.3).
- 9    For the execution of a **requeue** on an entry of a target protected object, after leaving the enclosing callable construct:
- 10
  - if the **requeue** is an internal **requeue** (that is, the **requeue** is back on an entry of the same protected object — see 9.5), the call is added to the queue of the named entry and the ongoing protected action continues (see 9.5.1);

- if the requeue is an external requeue (that is, the target protected object is not implicitly the same as the current object — see 9.5), a protected action is started on the target object and proceeds as for a normal entry call (see 9.5.3).

If the new entry named in the `requeue_statement` has formal parameters, then during the execution of the `accept_statement` or `entry_body` corresponding to the new entry, the formal parameters denote the same objects as did the corresponding formal parameters of the callable construct completed by the requeue. In any case, no parameters are specified in a `requeue_statement`; any parameter passing is implicit.

If the `requeue_statement` includes the reserved words **with abort** (it is a *requeue-with-abort*), then:

- if the original entry call has been aborted (see 9.8), then the requeue acts as an abort completion point for the call, and the call is cancelled and no requeue is performed;
- if the original entry call was timed (or conditional), then the original expiration time is the expiration time for the requeued call.

If the reserved words **with abort** do not appear, then the call remains protected against cancellation while queued as the result of the `requeue_statement`.

#### NOTES

31 A requeue is permitted from a single entry to an entry of an entry family, or vice-versa. The entry index, if any, plays no part in the subtype conformance check between the profiles of the two entries; an entry index is part of the `entry_name` for an entry of a family.

#### Examples

##### *Examples of requeue statements:*

```
requeue Request(Medium) with abort;
 -- requeue on a member of an entry family of the current task, see 9.1
requeue Flags(I).Seize;
 -- requeue on an entry of an array component, see 9.4
```

## 9.6 Delay Statements, Duration, and Time

A `delay_statement` is used to block further execution until a specified *expiration time* is reached. The expiration time can be specified either as a particular point in time (in a `delay_until_statement`), or in seconds from the current time (in a `delay_relative_statement`). The language-defined package `Calendar` provides definitions for a type `Time` and associated operations, including a function `Clock` that returns the current time.

#### Syntax

```
delay_statement ::= delay_until_statement | delay_relative_statement
delay_until_statement ::= delay until delay_expression;
delay_relative_statement ::= delay delay_expression;
```

#### Name Resolution Rules

The expected type for the `delay_expression` in a `delay_relative_statement` is the predefined type `Duration`. The `delay_expression` in a `delay_until_statement` is expected to be of any nonlimited type.

#### Legality Rules

There can be multiple time bases, each with a corresponding clock, and a corresponding *time type*. The type of the `delay_expression` in a `delay_until_statement` shall be a time type — either the type `Time`

defined in the language-defined package `Calendar` (see below), or some other implementation-defined time type (see D.8).

*Static Semantics*

- 7 There is a predefined fixed point type named `Duration`, declared in the visible part of package `Standard`; a value of type `Duration` is used to represent the length of an interval of time, expressed in seconds. The type `Duration` is not specific to a particular time base, but can be used with any time base.
- 8 A value of the type `Time` in package `Calendar`, or of some other implementation-defined time type, represents a time as reported by a corresponding clock.
- 9 The following language-defined library package exists:

```

10 package Ada.Calendar is
11 type Time is private;
12 subtype Year_Number is Integer range 1901 .. 2399;
13 subtype Month_Number is Integer range 1 .. 12;
14 subtype Day_Number is Integer range 1 .. 31;
15 subtype Day_Duration is Duration range 0.0 .. 86_400.0;
16 function Clock return Time;
17 function Year (Date : Time) return Year_Number;
18 function Month (Date : Time) return Month_Number;
19 function Day (Date : Time) return Day_Number;
20 function Seconds (Date : Time) return Day_Duration;
21 procedure Split (Date : in Time;
22 Year : out Year_Number;
23 Month : out Month_Number;
24 Day : out Day_Number;
25 Seconds : out Day_Duration);
26 function Time_Of (Year : Year_Number;
27 Month : Month_Number;
28 Day : Day_Number;
29 Seconds : Day_Duration := 0.0)
30 return Time;
31 function "+" (Left : Time; Right : Duration) return Time;
32 function "+" (Left : Duration; Right : Time) return Time;
33 function "-" (Left : Time; Right : Duration) return Time;
34 function "-" (Left : Time; Right : Time) return Duration;
35 function "<" (Left, Right : Time) return Boolean;
36 function "<=" (Left, Right : Time) return Boolean;
37 function ">" (Left, Right : Time) return Boolean;
38 function ">=" (Left, Right : Time) return Boolean;
39 Time_Error : exception;
40 private
41 ... -- not specified by the language
42 end Ada.Calendar;
```

*Dynamic Semantics*

- 20 For the execution of a `delay_statement`, the `delay_expression` is first evaluated. For a `delay_until_statement`, the expiration time for the delay is the value of the `delay_expression`, in the time base associated with the type of the expression. For a `delay_relative_statement`, the expiration time is defined as the current time, in the time base associated with relative delays, plus the value of the `delay_expression` converted to the type `Duration`, and then rounded up to the next clock tick. The time base associated with relative delays is as defined in D.9, “Delay Accuracy” or is implementation defined.

The task executing a `delay_statement` is blocked until the expiration time is reached, at which point it becomes ready again. If the expiration time has already passed, the task is not blocked.

If an attempt is made to *cancel* the `delay_statement` (as part of an `asynchronous_select` or `abort` — see 9.7.4 and 9.8), the `_statement` is cancelled if the expiration time has not yet passed, thereby completing the `delay_statement`.

The time base associated with the type `Time` of package `Calendar` is implementation defined. The function `Clock` of package `Calendar` returns a value representing the current time for this time base. The implementation-defined value of the named number `System.Tick` (see 13.7) is an approximation of the length of the real-time interval during which the value of `Calendar.Clock` remains constant.

The functions `Year`, `Month`, `Day`, and `Seconds` return the corresponding values for a given value of the type `Time`, as appropriate to an implementation-defined time zone; the procedure `Split` returns all four corresponding values. Conversely, the function `Time_Of` combines a year number, a month number, a day number, and a duration, into a value of type `Time`. The operators "+" and "-" for addition and subtraction of times and durations, and the relational operators for times, have the conventional meaning.

If `Time_Of` is called with a seconds value of `86_400.0`, the value returned is equal to the value of `Time_Of` for the next day with a seconds value of `0.0`. The value returned by the function `Seconds` or through the `Seconds` parameter of the procedure `Split` is always less than `86_400.0`.

The exception `Time_Error` is raised by the function `Time_Of` if the actual parameters do not form a proper date. This exception is also raised by the operators "+" and "-" if the result is not representable in the type `Time` or `Duration`, as appropriate. This exception is also raised by the functions `Year`, `Month`, `Day`, and `Seconds` and the procedure `Split` if the year number of the given date is outside of the range of the subtype `Year_Number`.

#### *Implementation Requirements*

The implementation of the type `Duration` shall allow representation of time intervals (both positive and negative) up to at least 86400 seconds (one day); `Duration'Small` shall not be greater than twenty milliseconds. The implementation of the type `Time` shall allow representation of all dates with year numbers in the range of `Year_Number`; it may allow representation of other dates as well (both earlier and later).

#### *Implementation Permissions*

An implementation may define additional time types (see D.8).

An implementation may raise `Time_Error` if the value of a `delay_expression` in a `delay_until_statement` of a `select_statement` represents a time more than 90 days past the current time. The actual limit, if any, is implementation-defined.

#### *Implementation Advice*

Whenever possible in an implementation, the value of `Duration'Small` should be no greater than 100 microseconds.

The time base for `delay_relative_statements` should be monotonic; it need not be the same time base as used for `Calendar.Clock`.

#### NOTES

32 A `delay_relative_statement` with a negative value of the `delay_expression` is equivalent to one with a zero value.

- 33    33 A `delay_statement` may be executed by the environment task; consequently `delay_statements` may be executed as part of the elaboration of a `library_item` or the execution of the main subprogram. Such statements delay the environment task (see 10.2).

34    34 A `delay_statement` is an abort completion point and a potentially blocking operation, even if the task is not actually blocked.

35    35 There is no necessary relationship between `System.Tick` (the resolution of the clock of package `Calendar`) and `Duration'Small` (the *small* of type `Duration`).

36    36 Additional requirements associated with `delay statements` are given in D.9, “Delay Accuracy”.

### *Examples*

37 Example of a relative delay statement:

38           **delay** 3.0;    -- delay 3.0 seconds

### 39 Example of a periodic task:

```

40 declare
 use Ada.Calendar;
 Next_Time : Time := Clock + Period;
 -- Period is a global constant of type Duration
begin
 loop -- repeated every Period seconds
 delay until Next_Time;
 ... -- perform some actions
 Next_Time := Next_Time + Period;
 end loop;
end;

```

### 9.6.1 Formatting, Time Zones, and other operations for Time

Static Semantics

- 1/2 The following language-defined library packages exist:

```

2/2 package Ada.Calendar.Time_Zones is
3/2 -- Time zone manipulation:
4/2 type Time_Offset is range -28*60 .. 28*60;
5/2 Unknown_Zone_Error : exception;
6/2 function UTC_Time_Offset (Date : Time := Clock) return Time_Offset;
7/2 end Ada.Calendar.Time_Zones;

8/2 package Ada.Calendar.Arithmetic is
9/2 -- Arithmetic on days:
10/2 type Day_Count is range
 -366*(1+Year_Number'Last - Year_Number'First)
 ..
 366*(1+Year_Number'Last - Year_Number'First);
11/2 subtype Leap_Seconds_Count is Integer range -2047 .. 2047;
12/2 procedure Difference (Left, Right : in Time;
 Days : out Day_Count;
 Seconds : out Duration;
 Leap_Seconds : out Leap_Seconds_Count);

13/2 function "+" (Left : Time; Right : Day_Count) return Time;
 function "+" (Left : Day_Count; Right : Time) return Time;
 function "-" (Left : Time; Right : Day_Count) return Time;
 function "--" (Left, Right : Time) return Day_Count;

```

```

end Ada.Calendar.Arithmetic; 14/2

with Ada.Calendar.Time_Zones; 15/2
package Ada.Calendar.Formatting is
 -- Day of the week: 16/2
 type Day_Name is (Monday, Tuesday, Wednesday, Thursday,
 Friday, Saturday, Sunday); 17/2
 function Day_of_Week (Date : Time) return Day_Name; 18/2
 -- Hours:Minutes:Seconds access: 19/2
 subtype Hour_Number is Natural range 0 .. 23; 20/2
 subtype Minute_Number is Natural range 0 .. 59;
 subtype Second_Number is Natural range 0 .. 59;
 subtype Second_Duration is Day_Duration range 0.0 .. 1.0;
 function Year (Date : Time;
 Time_Zone : Time_Zones.Time_Offset := 0)
 return Year_Number; 21/2
 function Month (Date : Time;
 Time_Zone : Time_Zones.Time_Offset := 0)
 return Month_Number; 22/2
 function Day (Date : Time;
 Time_Zone : Time_Zones.Time_Offset := 0)
 return Day_Number; 23/2
 function Hour (Date : Time;
 Time_Zone : Time_Zones.Time_Offset := 0)
 return Hour_Number; 24/2
 function Minute (Date : Time;
 Time_Zone : Time_Zones.Time_Offset := 0)
 return Minute_Number; 25/2
 function Second (Date : Time)
 return Second_Number; 26/2
 function Sub_Second (Date : Time)
 return Second_Duration; 27/2
 function Seconds_Of (Hour : Hour_Number;
 Minute : Minute_Number;
 Second : Second_Number := 0;
 Sub_Second : Second_Duration := 0.0)
 return Day_Duration; 28/2
 procedure Split (Seconds : in Day_Duration;
 Hour : out Hour_Number;
 Minute : out Minute_Number;
 Second : out Second_Number;
 Sub_Second : out Second_Duration); 29/2
 function Time_Of (Year : Year_Number;
 Month : Month_Number;
 Day : Day_Number;
 Hour : Hour_Number;
 Minute : Minute_Number;
 Second : Second_Number;
 Sub_Second : Second_Duration := 0.0;
 Leap_Second: Boolean := False;
 Time_Zone : Time_Zones.Time_Offset := 0)
 return Time; 30/2
 function Time_Of (Year : Year_Number;
 Month : Month_Number;
 Day : Day_Number;
 Seconds : Day_Duration := 0.0;
 Leap_Second: Boolean := False;
 Time_Zone : Time_Zones.Time_Offset := 0)
 return Time; 31/2

```

```

32/2 procedure Split (Date : in Time;
 Year : out Year_Number;
 Month : out Month_Number;
 Day : out Day_Number;
 Hour : out Hour_Number;
 Minute : out Minute_Number;
 Second : out Second_Number;
 Sub_Second: out Second_Duration;
 Time_Zone : in Time_Zones.Time_Offset := 0);

33/2 procedure Split (Date : in Time;
 Year : out Year_Number;
 Month : out Month_Number;
 Day : out Day_Number;
 Hour : out Hour_Number;
 Minute : out Minute_Number;
 Second : out Second_Number;
 Sub_Second: out Second_Duration;
 Leap_Second: out Boolean;
 Time_Zone : in Time_Zones.Time_Offset := 0);

34/2 procedure Split (Date : in Time;
 Year : out Year_Number;
 Month : out Month_Number;
 Day : out Day_Number;
 Seconds : out Day_Duration;
 Leap_Second: out Boolean;
 Time_Zone : in Time_Zones.Time_Offset := 0);

35/2 -- Simple image and value:
function Image (Date : Time;
 Include_Time_Fraction : Boolean := False;
 Time_Zone : Time_Zones.Time_Offset := 0) return String;

36/2 function Value (Date : String;
 Time_Zone : Time_Zones.Time_Offset := 0) return Time;

37/2 function Image (Elapsed_Time : Duration;
 Include_Time_Fraction : Boolean := False) return String;

38/2 function Value (Elapsed_Time : String) return Duration;

39/2 end Ada.Calendar.Formatting;

40/2 Type Time_Offset represents the number of minutes difference between the implementation-defined time
zone used by Calendar and another time zone.

41/2 function UTC_Time_Offset (Date : Time := Clock) return Time_Offset;

42/2 Returns, as a number of minutes, the difference between the implementation-defined time zone
of Calendar, and UTC time, at the time Date. If the time zone of the Calendar implementation is
unknown, then Unknown_Zone_Error is raised.

43/2 procedure Difference (Left, Right : in Time;
 Days : out Day_Count;
 Seconds : out Duration;
 Leap_Seconds : out Leap_Seconds_Count);

44/2 Returns the difference between Left and Right. Days is the number of days of difference,
Seconds is the remainder seconds of difference excluding leap seconds, and Leap_Seconds is the
number of leap seconds. If Left < Right, then Seconds <= 0.0, Days <= 0, and Leap_Seconds <=
0. Otherwise, all values are nonnegative. The absolute value of Seconds is always less than
86_400.0. For the returned values, if Days = 0, then Seconds + Duration(Leap_Seconds) =
Calendar."-" (Left, Right).

```

|                                                                                                                                      |      |
|--------------------------------------------------------------------------------------------------------------------------------------|------|
| <b>function</b> "+" (Left : Time; Right : Day_Count) <b>return</b> Time;                                                             | 45/2 |
| <b>function</b> "+" (Left : Day_Count; Right : Time) <b>return</b> Time;                                                             | 46/2 |
| Adds a number of days to a time value. Time_Error is raised if the result is not representable as a value of type Time.              |      |
| <b>function</b> "-" (Left : Time; Right : Day_Count) <b>return</b> Time;                                                             | 47/2 |
| Subtracts a number of days from a time value. Time_Error is raised if the result is not representable as a value of type Time.       |      |
| <b>function</b> "-" (Left, Right : Time) <b>return</b> Day_Count;                                                                    | 48/2 |
| Subtracts two time values, and returns the number of days between them. This is the same value that Difference would return in Days. |      |
| <b>function</b> Day_of_Week (Date : Time) <b>return</b> Day_Name;                                                                    | 49/2 |
| Returns the day of the week for Time. This is based on the Year, Month, and Day values of Time.                                      |      |
| <b>function</b> Year (Date : Time;<br>Time_Zone : Time_Zones.Time_Offset := 0)<br><b>return</b> Year_Number;                         | 50/2 |
| Returns the year for Date, as appropriate for the specified time zone offset.                                                        |      |
| <b>function</b> Month (Date : Time;<br>Time_Zone : Time_Zones.Time_Offset := 0)<br><b>return</b> Month_Number;                       | 51/2 |
| Returns the month for Date, as appropriate for the specified time zone offset.                                                       |      |
| <b>function</b> Day (Date : Time;<br>Time_Zone : Time_Zones.Time_Offset := 0)<br><b>return</b> Day_Number;                           | 52/2 |
| Returns the day number for Date, as appropriate for the specified time zone offset.                                                  |      |
| <b>function</b> Hour (Date : Time;<br>Time_Zone : Time_Zones.Time_Offset := 0)<br><b>return</b> Hour_Number;                         | 53/2 |
| Returns the hour for Date, as appropriate for the specified time zone offset.                                                        |      |
| <b>function</b> Minute (Date : Time;<br>Time_Zone : Time_Zones.Time_Offset := 0)<br><b>return</b> Minute_Number;                     | 54/2 |
| Returns the minute within the hour for Date, as appropriate for the specified time zone offset.                                      |      |
| <b>function</b> Second (Date : Time)<br><b>return</b> Second_Number;                                                                 | 55/2 |
| Returns the second within the hour and minute for Date.                                                                              |      |
| <b>function</b> Sub_Second (Date : Time)<br><b>return</b> Second_Duration;                                                           | 56/2 |
| Returns the fraction of second for Date (this has the same accuracy as Day_Duration). The value returned is always less than 1.0.    |      |

67/2       **function** Seconds\_Of (Hour   : Hour\_Number;  
                  Minute : Minute\_Number;  
                  Second : Second\_Number := 0;  
                  Sub\_Second : Second\_Duration := 0.0)  
       **return** Day\_Duration;

68/2       Returns a Day\_Duration value for the combination of the given Hour, Minute, Second, and Sub\_Second. This value can be used in Calendar.Time.Of as well as the argument to Calendar."+" and Calendar.".−". If Seconds.Of is called with a Sub\_Second value of 1.0, the value returned is equal to the value of Seconds.Of for the next second with a Sub\_Second value of 0.0.

69/2       **procedure** Split (Seconds   : **in** Day\_Duration;  
                  Hour      : **out** Hour\_Number;  
                  Minute   : **out** Minute\_Number;  
                  Second   : **out** Second\_Number;  
                  Sub\_Second : **out** Second\_Duration);

70/2       Splits Seconds into Hour, Minute, Second and Sub\_Second in such a way that the resulting values all belong to their respective subtypes. The value returned in the Sub\_Second parameter is always less than 1.0.

71/2       **function** Time.Of (Year       : Year\_Number;  
                  Month     : Month\_Number;  
                  Day       : Day\_Number;  
                  Hour      : Hour\_Number;  
                  Minute   : Minute\_Number;  
                  Second   : Second\_Number;  
                  Sub\_Second : Second\_Duration := 0.0;  
                  Leap\_Second: Boolean := False;  
                  Time\_Zone : Time\_Zones.Time\_Offset := 0)  
       **return** Time;

72/2       If Leap\_Second is False, returns a Time built from the date and time values, relative to the specified time zone offset. If Leap\_Second is True, returns the Time that represents the time within the leap second that is one second later than the time specified by the other parameters. Time\_Error is raised if the parameters do not form a proper date or time. If Time.Of is called with a Sub\_Second value of 1.0, the value returned is equal to the value of Time.Of for the next second with a Sub\_Second value of 0.0.

73/2       **function** Time.Of (Year       : Year\_Number;  
                  Month     : Month\_Number;  
                  Day       : Day\_Number;  
                  Seconds   : Day\_Duration := 0.0;  
                  Leap\_Second: Boolean := False;  
                  Time\_Zone : Time\_Zones.Time\_Offset := 0)  
       **return** Time;

74/2       If Leap\_Second is False, returns a Time built from the date and time values, relative to the specified time zone offset. If Leap\_Second is True, returns the Time that represents the time within the leap second that is one second later than the time specified by the other parameters. Time\_Error is raised if the parameters do not form a proper date or time. If Time.Of is called with a Seconds value of 86\_400.0, the value returned is equal to the value of Time.Of for the next day with a Seconds value of 0.0.

```
procedure Split (Date : in Time;
 Year : out Year_Number;
 Month : out Month_Number;
 Day : out Day_Number;
 Hour : out Hour_Number;
 Minute : out Minute_Number;
 Second : out Second_Number;
 Sub_Second : out Second_Duration;
 Leap_Second: out Boolean;
 Time_Zone : in Time_Zones.Time_Offset := 0);
```

75/2

If Date does not represent a time within a leap second, splits Date into its constituent parts (Year, Month, Day, Hour, Minute, Second, Sub\_Second), relative to the specified time zone offset, and sets Leap\_Second to False. If Date represents a time within a leap second, set the constituent parts to values corresponding to a time one second earlier than that given by Date, relative to the specified time zone offset, and sets Leap\_Seconds to True. The value returned in the Sub\_Second parameter is always less than 1.0.

76/2

```
procedure Split (Date : in Time;
 Year : out Year_Number;
 Month : out Month_Number;
 Day : out Day_Number;
 Hour : out Hour_Number;
 Minute : out Minute_Number;
 Second : out Second_Number;
 Sub_Second : out Second_Duration;
 Time_Zone : in Time_Zones.Time_Offset := 0);
```

77/2

Splits Date into its constituent parts (Year, Month, Day, Hour, Minute, Second, Sub\_Second), relative to the specified time zone offset. The value returned in the Sub\_Second parameter is always less than 1.0.

78/2

```
procedure Split (Date : in Time;
 Year : out Year_Number;
 Month : out Month_Number;
 Day : out Day_Number;
 Seconds : out Day_Duration;
 Leap_Second: out Boolean;
 Time_Zone : in Time_Zones.Time_Offset := 0);
```

79/2

If Date does not represent a time within a leap second, splits Date into its constituent parts (Year, Month, Day, Seconds), relative to the specified time zone offset, and sets Leap\_Second to False. If Date represents a time within a leap second, set the constituent parts to values corresponding to a time one second earlier than that given by Date, relative to the specified time zone offset, and sets Leap\_Seconds to True. The value returned in the Seconds parameter is always less than 86\_400.0.

80/2

```
function Image (Date : Time;
 Include_Time_Fraction : Boolean := False;
 Time_Zone : Time_Zones.Time_Offset := 0) return String;
```

81/2

Returns a string form of the Date relative to the given Time\_Zone. The format is "Year-Month-Day Hour:Minute:Second", where the Year is a 4-digit value, and all others are 2-digit values, of the functions defined in Calendar and Calendar.Formatting, including a leading zero, if needed. The separators between the values are a minus, another minus, a colon, and a single space between the Day and Hour. If Include\_Time\_Fraction is True, the integer part of Sub\_Seconds\*100 is suffixed to the string as a point followed by a 2-digit value.

82/2

```

83/2 function Value (Date : String;
 Time_Zone : Time_Zones.Time_Offset := 0) return Time;
84/2 Returns a Time value for the image given as Date, relative to the given time zone.
 Constraint_Error is raised if the string is not formatted as described for Image, or the function
 cannot interpret the given string as a Time value.

85/2 function Image (Elapsed_Time : Duration;
 Include_Time_Fraction : Boolean := False) return String;
86/2 Returns a string form of the Elapsed_Time. The format is "Hour:Minute:Second", where all
 values are 2-digit values, including a leading zero, if needed. The separators between the values
 are colons. If Include_Time_Fraction is True, the integer part of Sub_Seconds*100 is suffixed to
 the string as a point followed by a 2-digit value. If Elapsed_Time < 0.0, the result is Image (abs
 Elapsed_Time, Include_Time_Fraction) prefixed with a minus sign. If abs Elapsed_Time
 represents 100 hours or more, the result is implementation-defined.

87/2 function Value (Elapsed_Time : String) return Duration;
88/2 Returns a Duration value for the image given as Elapsed_Time. Constraint_Error is raised if the
 string is not formatted as described for Image, or the function cannot interpret the given string as
 a Duration value.

```

*Implementation Advice*

89/2 An implementation should support leap seconds if the target system supports them. If leap seconds are not supported, Difference should return zero for Leap\_Seconds, Split should return False for Leap\_Second, and Time\_Of should raise Time\_Error if Leap\_Second is True.

NOTES

90/2 37 The implementation-defined time zone of package Calendar may, but need not, be the local time zone. UTC\_Time\_Offset always returns the difference relative to the implementation-defined time zone of package Calendar. If UTC\_Time\_Offset does not raise Unknown\_Zone\_Error, UTC time can be safely calculated (within the accuracy of the underlying time-base).

91/2 38 Calling Split on the results of subtracting Duration(UTC\_Time\_Offset\*60) from Clock provides the components (hours, minutes, and so on) of the UTC time. In the United States, for example, UTC\_Time\_Offset will generally be negative.

## 9.7 Select Statements

1 There are four forms of the `select_statement`. One form provides a selective wait for one or more `select_alternatives`. Two provide timed and conditional entry calls. The fourth provides asynchronous transfer of control.

*Syntax*

```

2 select_statement ::=
 selective_accept
 | timed_entry_call
 | conditional_entry_call
 | asynchronous_select

```

*Examples**Example of a select statement:*

```

select
 accept Driver_Awake_Signal;
or
 delay 30.0*Seconds;
 Stop_The_Train;
end select;
```

3  
4**9.7.1 Selective Accept**

This form of the `select_statement` allows a combination of waiting for, and selecting from, one or more alternatives. The selection may depend on conditions associated with each alternative of the `selective_accept`.

1

*Syntax*

```

selective_accept ::= 2
 select
 [guard]
 select_alternative
 { or
 [guard]
 select_alternative }
 [else
 sequence_of_statements]
 end select;

guard ::= when condition => 3
select_alternative ::= 4
 accept_alternative
 | delay_alternative
 | terminate_alternative
accept_alternative ::= 5
 accept_statement [sequence_of_statements]
delay_alternative ::= 6
 delay_statement [sequence_of_statements]
terminate_alternative ::= terminate; 7
```

A `selective_accept` shall contain at least one `accept_alternative`. In addition, it can contain:

- a `terminate_alternative` (only one); or
- one or more `delay_alternatives`; or
- an `else part` (the reserved word `else` followed by a `sequence_of_statements`).

These three possibilities are mutually exclusive.

8  
9  
10  
11  
12  
13*Legality Rules*

If a `selective_accept` contains more than one `delay_alternative`, then all shall be `delay_relative_-statements`, or all shall be `delay_until_statements` for the same time type.

14

*Dynamic Semantics*

A `select_alternative` is said to be *open* if it is not immediately preceded by a `guard`, or if the condition of its `guard` evaluates to True. It is said to be *closed* otherwise.

14

- 15 For the execution of a `selective_accept`, any guard conditions are evaluated; open alternatives are thus determined. For an open `delay_alternative`, the `delay_expression` is also evaluated. Similarly, for an open `accept_alternative` for an entry of a family, the `entry_index` is also evaluated. These evaluations are performed in an arbitrary order, except that a `delay_expression` or `entry_index` is not evaluated until after evaluating the corresponding condition, if any. Selection and execution of one open alternative, or of the `else` part, then completes the execution of the `selective_accept`; the rules for this selection are described below.
- 16 Open `accept_alternatives` are first considered. Selection of one such alternative takes place immediately if the corresponding entry already has queued calls. If several alternatives can thus be selected, one of them is selected according to the entry queuing policy in effect (see 9.5.3 and D.4). When such an alternative is selected, the selected call is removed from its entry queue and the `handled_sequence_of_statements` (if any) of the corresponding `accept_statement` is executed; after the rendezvous completes any subsequent `sequence_of_statements` of the alternative is executed. If no selection is immediately possible (in the above sense) and there is no `else` part, the task blocks until an open alternative can be selected.
- 17 Selection of the other forms of alternative or of an `else` part is performed as follows:
- 18 • An open `delay_alternative` is selected when its expiration time is reached if no `accept_alternative` or other `delay_alternative` can be selected prior to the expiration time. If several `delay_alternatives` have this same expiration time, one of them is selected according to the queuing policy in effect (see D.4); the default queuing policy chooses arbitrarily among the `delay_alternatives` whose expiration time has passed.
  - 19 • The `else` part is selected and its `sequence_of_statements` is executed if no `accept_alternative` can immediately be selected; in particular, if all alternatives are closed.
  - 20 • An open `terminate_alternative` is selected if the conditions stated at the end of clause 9.3 are satisfied.
- 21 The exception `Program_Error` is raised if all alternatives are closed and there is no `else` part.

## NOTES

22 39 A `selective_accept` is allowed to have several open `delay_alternatives`. A `selective_accept` is allowed to have several open `accept_alternatives` for the same entry.

*Examples*

- 23 Example of a task body with a `selective accept`:

```
24 task body Server is
 Current_Work_Item : Work_Item;
 begin
 loop
 select
 accept Next_Work_Item(WI : in Work_Item) do
 Current_Work_Item := WI;
 end;
 Process_Work_Item(Current_Work_Item);
 or
 accept Shut_Down;
 exit; -- Premature shut down requested
 or
 terminate; -- Normal shutdown at end of scope
 end select;
 end loop;
 end Server;
```

## 9.7.2 Timed Entry Calls

A `timed_entry_call` issues an entry call that is cancelled if the call (or a requeue-with-abort of the call) is not selected before the expiration time is reached. A procedure call may appear rather than an entry call for cases where the procedure might be implemented by an entry.

*Syntax*

```

timed_entry_call ::= 2
 select
 entry_call_alternative
 or
 delay_alternative
 end select;

entry_call_alternative ::= 3/2
 procedure_or_entry_call [sequence_of_statements]

procedure_or_entry_call ::= 3.1/2
 procedure_call_statement | entry_call_statement

```

*Legality Rules*

If a `procedure_call_statement` is used for a `procedure_or_entry_call`, the `procedure_name` or `procedure_prefix` of the `procedure_call_statement` shall statically denote an entry renamed as a procedure or (a view of) a primitive subprogram of a limited interface whose first parameter is a controlling parameter (see 3.9.2).

*Static Semantics*

If a `procedure_call_statement` is used for a `procedure_or_entry_call`, and the procedure is implemented by an entry, then the `procedure_name`, or `procedure_prefix` and possibly the first parameter of the `procedure_call_statement`, determine the target object of the call and the entry to be called.

*Dynamic Semantics*

For the execution of a `timed_entry_call`, the `entry_name`, `procedure_name`, or `procedure_prefix`, and any actual parameters are evaluated, as for a simple entry call (see 9.5.3) or procedure call (see 6.4). The expiration time (see 9.6) for the call is determined by evaluating the `delay_expression` of the `delay_alternative`. If the call is an entry call or a call on a procedure implemented by an entry, the entry call is then issued. Otherwise, the call proceeds as described in 6.4 for a procedure call, followed by the `sequence_of_statements` of the `entry_call_alternative`; the `sequence_of_statements` of the `delay_alternative` is ignored.

If the call is queued (including due to a requeue-with-abort), and not selected before the expiration time is reached, an attempt to cancel the call is made. If the call completes due to the cancellation, the optional `sequence_of_statements` of the `delay_alternative` is executed; if the entry call completes normally, the optional `sequence_of_statements` of the `entry_call_alternative` is executed.

*Examples*

6    Example of a timed entry call:

```
7 select
 Controller.Request (Medium) (Some_Item);
or
 delay 45.0;
 -- controller too busy, try something else
end select;
```

### 9.7.3 Conditional Entry Calls

1/2 A conditional\_entry\_call issues an entry call that is then cancelled if it is not selected immediately (or if a requeue-with-abort of the call is not selected immediately). A procedure call may appear rather than an entry call for cases where the procedure might be implemented by an entry.

*Syntax*

```
2 conditional_entry_call ::=
select
entry_call_alternative
else
sequence_of_statements
end select;
```

*Dynamic Semantics*

3    The execution of a conditional\_entry\_call is defined to be equivalent to the execution of a timed\_entry\_call with a delay\_alternative specifying an immediate expiration time and the same sequence\_of\_statements as given after the reserved word else.

## NOTES

4    40 A conditional\_entry\_call may briefly increase the Count attribute of the entry, even if the conditional call is not selected.

*Examples*

5    Example of a conditional entry call:

```
6 procedure Spin(R : in Resource) is
begin
 loop
 select
 R.Seize;
 return;
 else
 null; -- busy waiting
 end select;
 end loop;
end;
```

## 9.7.4 Asynchronous Transfer of Control

An asynchronous `select_statement` provides asynchronous transfer of control upon completion of an entry call or the expiration of a delay.

### Syntax

```
asynchronous_select ::= 2
 select
 triggering_alternative
 then abort
 abortable_part
 end select;

triggering_alternative ::= triggering_statement [sequence_of_statements] 3
triggering_statement ::= procedure_or_entry_call | delay_statement 4/2
abortable_part ::= sequence_of_statements 5
```

### Dynamic Semantics

For the execution of an `asynchronous_select` whose `triggering_statement` is a `procedure_or_entry_call`, the `entry_name`, `procedure_name`, or `procedure_prefix`, and actual parameters are evaluated as for a simple entry call (see 9.5.3) or procedure call (see 6.4). If the call is an entry call or a call on a procedure implemented by an entry, the entry call is issued. If the entry call is queued (or requeued-with-abort), then the `abortable_part` is executed. If the entry call is selected immediately, and never requeued-with-abort, then the `abortable_part` is never started. If the call is on a procedure that is not implemented by an entry, the call proceeds as described in 6.4, followed by the `sequence_of_statements` of the `triggering_alternative`; the `abortable_part` is never started.

For the execution of an `asynchronous_select` whose `triggering_statement` is a `delay_statement`, the `delay_expression` is evaluated and the expiration time is determined, as for a normal `delay_statement`. If the expiration time has not already passed, the `abortable_part` is executed.

If the `abortable_part` completes and is left prior to completion of the `triggering_statement`, an attempt to cancel the `triggering_statement` is made. If the attempt to cancel succeeds (see 9.5.3 and 9.6), the `asynchronous_select` is complete.

If the `triggering_statement` completes other than due to cancellation, the `abortable_part` is aborted (if started but not yet completed — see 9.8). If the `triggering_statement` completes normally, the optional `sequence_of_statements` of the `triggering_alternative` is executed after the `abortable_part` is left.

### Examples

*Example of a main command loop for a command interpreter:*

```
loop 10
 select
 Terminal.Wait_For_Interrupt;
 Put_Line("Interrupted");
 then abort 11
 -- This will be abandoned upon terminal interrupt
 Put_Line("-> ");
 Get_Line(Command, Last);
 Process_Command(Command(1..Last));
 end select;
end loop;
```

*Example of a time-limited calculation:*

```

12 select
13 delay 5.0;
 Put_Line("Calculation does not converge");
14 then abort
15 -- This calculation should finish in 5.0 seconds;
16 -- if not, it is assumed to diverge.
17 Horribly_Complicated_Recursive_Function(X, Y);
18 end select;
```

## 9.8 Abort of a Task - Abort of a Sequence of Statements

- 1 An `abort_statement` causes one or more tasks to become abnormal, thus preventing any further interaction with such tasks. The completion of the `triggering_statement` of an `asynchronous_select` causes a `sequence_of_statements` to be aborted.

*Syntax*

2 `abort_statement ::= abort task_name {, task_name};`

*Name Resolution Rules*

- 3 Each `task_name` is expected to be of any task type; they need not all be of the same task type.

*Dynamic Semantics*

- 4 For the execution of an `abort_statement`, the given `task_names` are evaluated in an arbitrary order. Each named task is then *aborted*, which consists of making the task *abnormal* and aborting the execution of the corresponding `task_body`, unless it is already completed.
- 5 When the execution of a construct is *aborted* (including that of a `task_body` or of a `sequence_of_statements`), the execution of every construct included within the aborted execution is also aborted, except for executions included within the execution of an *abort-deferred* operation; the execution of an *abort-deferred* operation continues to completion without being affected by the abort; the following are the *abort-deferred* operations:

- 6 • a protected action;
- 7 • waiting for an entry call to complete (after having initiated the attempt to cancel it — see below);
- 8 • waiting for the termination of dependent tasks;
- 9 • the execution of an `Initialize` procedure as the last step of the default initialization of a controlled object;
- 10 • the execution of a `Finalize` procedure as part of the finalization of a controlled object;
- 11 • an assignment operation to an object with a controlled part.

12 The last three of these are discussed further in 7.6.

13 When a master is aborted, all tasks that depend on that master are aborted.

14 The order in which tasks become abnormal as the result of an `abort_statement` or the abort of a `sequence_of_statements` is not specified by the language.

15 If the execution of an entry call is aborted, an immediate attempt is made to cancel the entry call (see 9.5.3). If the execution of a construct is aborted at a time when the execution is blocked, other than for an entry call, at a point that is outside the execution of an *abort-deferred* operation, then the execution of the

construct completes immediately. For an abort due to an `abort_statement`, these immediate effects occur before the execution of the `abort_statement` completes. Other than for these immediate cases, the execution of a construct that is aborted does not necessarily complete before the `abort_statement` completes. However, the execution of the aborted construct completes no later than its next *abort completion point* (if any) that occurs outside of an abort-deferred operation; the following are abort completion points for an execution:

- the point where the execution initiates the activation of another task; 16
- the end of the activation of a task; 17
- the start or end of the execution of an entry call, `accept_statement`, `delay_statement`, or 18  
`abort_statement`;
- the start of the execution of a `select_statement`, or of the `sequence_of_statements` of an 19  
`exception_handler`.

#### *Bounded (Run-Time) Errors*

An attempt to execute an `asynchronous_select` as part of the execution of an abort-deferred operation is a 20 bounded error. Similarly, an attempt to create a task that depends on a master that is included entirely within the execution of an abort-deferred operation is a bounded error. In both cases, `Program_Error` is raised if the error is detected by the implementation; otherwise the operations proceed as they would outside an abort-deferred operation, except that an abort of the `abortable_part` or the created task might or might not have an effect.

#### *Erroneous Execution*

If an assignment operation completes prematurely due to an abort, the assignment is said to be *disrupted*; 21 the target of the assignment or its parts can become abnormal, and certain subsequent uses of the object can be erroneous, as explained in 13.9.1.

#### NOTES

- 41 An `abort_statement` should be used only in situations requiring unconditional termination. 22
- 42 A task is allowed to abort any task it can name, including itself. 23
- 43 Additional requirements associated with abort are given in D.6, “Preemptive Abort”. 24

## 9.9 Task and Entry Attributes

#### *Dynamic Semantics*

For a prefix T that is of a task type (after any implicit dereference), the following attributes are defined: 1

T'Callable      Yields the value True when the task denoted by T is *callable*, and False otherwise; a task is 2 callable unless it is completed or abnormal. The value of this attribute is of the predefined type Boolean.

T'Terminated    Yields the value True if the task denoted by T is terminated, and False otherwise. The value 3 of this attribute is of the predefined type Boolean.

For a prefix E that denotes an entry of a task or protected unit, the following attribute is defined. This 4 attribute is only allowed within the body of the task or protected unit, but excluding, in the case of an entry of a task unit, within any program unit that is, itself, inner to the body of the task unit.

E'Count        Yields the number of calls presently queued on the entry E of the current instance of the 5 unit. The value of this attribute is of the type `universal_integer`.

## NOTES

- 6    44 For the Count attribute, the entry can be either a single entry or an entry of a family. The name of the entry or entry family can be either a direct\_name or an expanded name.
- 7    45 Within task units, algorithms interrogating the attribute E'Count should take precautions to allow for the increase of the value of this attribute for incoming entry calls, and its decrease, for example with `timed_entry_calls`. Also, a `conditional_entry_call` may briefly increase this value, even if the conditional call is not accepted.
- 8    46 Within protected units, algorithms interrogating the attribute E'Count in the `entry_barrier` for the entry E should take precautions to allow for the evaluation of the condition of the barrier both before and after queuing a given caller.

## 9.10 Shared Variables

### *Static Semantics*

- 1    If two different objects, including nonoverlapping parts of the same object, are *independently addressable*, they can be manipulated concurrently by two different tasks without synchronization. Normally, any two nonoverlapping objects are independently addressable. However, if packing, record layout, or Component\_Size is specified for a given composite object, then it is implementation defined whether or not two nonoverlapping parts of that composite object are independently addressable.

### *Dynamic Semantics*

- 2    Separate tasks normally proceed independently and concurrently with one another. However, task interactions can be used to synchronize the actions of two or more tasks to allow, for example, meaningful communication by the direct updating and reading of variables shared between the tasks. The actions of two different tasks are synchronized in this sense when an action of one task *signals* an action of the other task; an action A1 is defined to signal an action A2 under the following circumstances:
- 3    • If A1 and A2 are part of the execution of the same task, and the language rules require A1 to be performed before A2;
  - 4    • If A1 is the action of an activator that initiates the activation of a task, and A2 is part of the execution of the task that is activated;
  - 5    • If A1 is part of the activation of a task, and A2 is the action of waiting for completion of the activation;
  - 6    • If A1 is part of the execution of a task, and A2 is the action of waiting for the termination of the task;
  - 6.1/1    • If A1 is the termination of a task T, and A2 is either the evaluation of the expression T'Terminated or a call to `Ada.Task_Identification.Is_Terminated` with an actual parameter that identifies T (see C.7.1);
  - 7    • If A1 is the action of issuing an entry call, and A2 is part of the corresponding execution of the appropriate `entry_body` or `accept_statement`.
  - 8    • If A1 is part of the execution of an `accept_statement` or `entry_body`, and A2 is the action of returning from the corresponding entry call;
  - 9    • If A1 is part of the execution of a protected procedure body or `entry_body` for a given protected object, and A2 is part of a later execution of an `entry_body` for the same protected object;
  - 10    • If A1 signals some action that in turn signals A2.

*Erroneous Execution*

Given an action of assigning to an object, and an action of reading or updating a part of the same object (or of a neighboring object if the two are not independently addressable), then the execution of the actions is erroneous unless the actions are *sequential*. Two actions are sequential if one of the following is true:

- One action signals the other;
- Both actions occur as part of the execution of the same task;
- Both actions occur as part of protected actions on the same protected object, and at most one of the actions is part of a call on a protected function of the protected object.

A pragma Atomic or Atomic\_Components may also be used to ensure that certain reads and updates are sequential — see C.6.

## 9.11 Example of Tasking and Synchronization

*Examples*

The following example defines a buffer protected object to smooth variations between the speed of output of a producing task and the speed of input of some consuming task. For instance, the producing task might have the following structure:

```
task Producer;
task body Producer is
 Person : Person_Name; -- see 3.10.1
begin
 loop
 ... -- simulate arrival of the next customer
 Buffer.Append_Wait(Person);
 exit when Person = null;
 end loop;
end Producer;
```

and the consuming task might have the following structure:

```
task Consumer;
task body Consumer is
 Person : Person_Name;
begin
 loop
 Buffer.Remove_First_Wait(Person);
 exit when Person = null;
 ... -- simulate serving a customer
 end loop;
end Consumer;
```

The buffer object contains an internal array of person names managed in a round-robin fashion. The array has two indices, an In\_Index denoting the index for the next input person name and an Out\_Index denoting the index for the next output person name.

The Buffer is defined as an extension of the Synchronized\_Queue interface (see 3.9.4), and as such promises to implement the abstraction defined by that interface. By doing so, the Buffer can be passed to the Transfer class-wide operation defined for objects of a type covered by Queue'Class.

```

8/2 protected Buffer is new Synchronized_Queue with -- see 3.9.4
 entry Append_Wait(Person : in Person_Name);
 entry Remove_First_Wait(Person : out Person_Name);
 function Cur_Count return Natural;
 function Max_Count return Natural;
 procedure Append(Person : in Person_Name);
 procedure Remove_First(Person : out Person_Name);
private
 Pool : Person_Name_Array(1 .. 100);
 Count : Natural := 0;
 In_Index, Out_Index : Positive := 1;
end Buffer;

9/2 protected body Buffer is
 entry Append_Wait(Person : in Person_Name)
 when Count < Pool'Length is
 begin
 Append(Person);
 end Append_Wait;

9.1/2 procedure Append(Person : in Person_Name) is
begin
 if Count = Pool'Length then
 raise Queue_Error with "Buffer Full"; -- see 11.3
 end if;
 Pool(In_Index) := Person;
 In_Index := (In_Index mod Pool'Length) + 1;
 Count := Count + 1;
end Append;

10/2 entry Remove_First_Wait(Person : out Person_Name)
 when Count > 0 is
 begin
 Remove_First(Person);
 end Remove_First_Wait;

11/2 procedure Remove_First(Person : out Person_Name) is
begin
 if Count = 0 then
 raise Queue_Error with "Buffer Empty"; -- see 11.3
 end if;
 Person := Pool(Out_Index);
 Out_Index := (Out_Index mod Pool'Length) + 1;
 Count := Count - 1;
end Remove_First;

12/2 function Cur_Count return Natural is
begin
 return Buffer.Count;
end Cur_Count;

13/2 function Max_Count return Natural is
begin
 return Pool'Length;
end Max_Count;
end Buffer;

```

## Section 10: Program Structure and Compilation Issues

The overall structure of programs and the facilities for separate compilation are described in this section. A *program* is a set of *partitions*, each of which may execute in a separate address space, possibly on a separate computer.

As explained below, a partition is constructed from *library units*. Syntactically, the declaration of a library unit is a *library\_item*, as is the body of a library unit. An implementation may support a concept of a *program library* (or simply, a “library”), which contains *library\_items* and their subunits. Library units may be organized into a hierarchy of children, grandchildren, and so on.

This section has two clauses: 10.1, “Separate Compilation” discusses compile-time issues related to separate compilation. 10.2, “Program Execution” discusses issues related to what is traditionally known as “link time” and “run time” — building and executing partitions.

### 10.1 Separate Compilation

A *program unit* is either a package, a task unit, a protected unit, a protected entry, a generic unit, or an explicitly declared subprogram other than an enumeration literal. Certain kinds of program units can be separately compiled. Alternatively, they can appear physically nested within other program units.

The text of a program can be submitted to the compiler in one or more compilations. Each *compilation* is a succession of *compilation\_units*. A *compilation\_unit* contains either the declaration, the body, or a renaming of a program unit. The representation for a *compilation* is implementation-defined.

A library unit is a separately compiled program unit, and is always a package, subprogram, or generic unit. Library units may have other (logically nested) library units as children, and may have other program units physically nested within them. A root library unit, together with its children and grandchildren and so on, form a *subsystem*.

#### *Implementation Permissions*

An implementation may impose implementation-defined restrictions on compilations that contain multiple *compilation\_units*.

#### 10.1.1 Compilation Units - Library Units

A *library\_item* is a compilation unit that is the declaration, body, or renaming of a library unit. Each library unit (except Standard) has a *parent unit*, which is a library package or generic library package. A library unit is a *child* of its parent unit. The *root* library units are the children of the predefined library package Standard.

#### *Syntax*

```

compilation ::= {compilation_unit}
compilation_unit ::=
 context_clause library_item
 | context_clause subunit
library_item ::= [private] library_unit_declaration
 | library_unit_body
 | [private] library_unit_renaming_declaration

```

```

5 library_unit_declaration ::=
 subprogram_declaration | package_declaration
 | generic_declaration | generic_instantiation
6 library_unit_renaming_declaration ::=
 package_renaming_declaration
 | generic_renaming_declaration
 | subprogram_renaming_declaration
7 library_unit_body ::= subprogram_body | package_body
8 parent_unit_name ::= name

```

8.1/2 An overriding\_indicator is not allowed in a subprogram\_declaration, generic\_instantiation, or subprogram\_renaming\_declaration that declares a library unit.

9 A *library unit* is a program unit that is declared by a *library\_item*. When a program unit is a library unit, the prefix “library” is used to refer to it (or “generic library” if generic), as well as to its declaration and body, as in “library procedure”, “library package\_body”, or “generic library package”. The term *compilation unit* is used to refer to a *compilation\_unit*. When the meaning is clear from context, the term is also used to refer to the *library\_item* of a *compilation\_unit* or to the *proper\_body* of a *subunit* (that is, the *compilation\_unit* without the *context\_clause* and the **separate** (*parent\_unit\_name*)).

10 The *parent declaration* of a *library\_item* (and of the library unit) is the declaration denoted by the *parent\_unit\_name*, if any, of the *defining\_program\_unit\_name* of the *library\_item*. If there is no *parent\_unit\_name*, the parent declaration is the declaration of Standard, the *library\_item* is a *root library\_item*, and the library unit (renaming) is a *root library unit* (renaming). The declaration and body of Standard itself have no parent declaration. The *parent unit* of a *library\_item* or library unit is the library unit declared by its parent declaration.

11 The children of a library unit occur immediately within the declarative region of the declaration of the library unit. The *ancestors* of a library unit are itself, its parent, its parent's parent, and so on. (Standard is an ancestor of every library unit.) The *descendant* relation is the inverse of the ancestor relation.

12 A *library\_unit\_declaration* or a *library\_unit\_renaming\_declaration* is *private* if the declaration is immediately preceded by the reserved word **private**; it is otherwise *public*. A library unit is private or public according to its declaration. The *public descendants* of a library unit are the library unit itself, and the public descendants of its public children. Its other descendants are *private descendants*.

12.1/2 For each library package\_declaration in the environment, there is an implicit declaration of a *limited view* of that library package. The limited view of a package contains:

- 12.2/2 • For each nested package\_declaration, a declaration of the limited view of that package, with the same *defining\_program\_unit\_name*.

- 12.3/2 • For each type\_declaration in the visible part, an incomplete view of the type; if the type\_declaration is tagged, then the view is a tagged incomplete view.

12.4/2 The limited view of a library package\_declaration is private if that library package\_declaration is immediately preceded by the reserved word **private**.

12.5/2 There is no syntax for declaring limited views of packages, because they are always implicit. The implicit declaration of a limited view of a library package is not the declaration of a library unit (the library package\_declaration is); nonetheless, it is a *library\_item*. The implicit declaration of the limited view of a library package forms an (implicit) compilation unit whose *context\_clause* is empty.

12.6/2 A library package\_declaration is the completion of the declaration of its limited view.

*Legality Rules*

The parent unit of a `library_item` shall be a library package or generic library package.

13

If a `defining_program_unit_name` of a given declaration or body has a `parent_unit_name`, then the given declaration or body shall be a `library_item`. The body of a program unit shall be a `library_item` if and only if the declaration of the program unit is a `library_item`. In a `library_unit_renaming_declaration`, the (old) name shall denote a `library_item`.

14

A `parent_unit_name` (which can be used within a `defining_program_unit_name` of a `library_item` and in the `separate` clause of a subunit), and each of its prefixes, shall not denote a `renaming_declaration`. On the other hand, a name that denotes a `library_unit_renaming_declaration` is allowed in a `nonlimited_with_clause` and other places where the name of a library unit is allowed.

15/2

If a library package is an instance of a generic package, then every child of the library package shall either be itself an instance or be a renaming of a library unit.

16

A child of a generic library package shall either be itself a generic unit or be a renaming of some other child of the same generic unit. The renaming of a child of a generic package shall occur only within the declarative region of the generic package.

17

A child of a parent generic package shall be instantiated or renamed only within the declarative region of the parent generic.

18

For each child  $C$  of some parent generic package  $P$ , there is a corresponding declaration  $C$  nested immediately within each instance of  $P$ . For the purposes of this rule, if a child  $C$  itself has a child  $D$ , each corresponding declaration for  $C$  has a corresponding child  $D$ . The corresponding declaration for a child within an instance is visible only within the scope of a `with_clause` that mentions the (original) child generic unit.

19/2

A library subprogram shall not override a primitive subprogram.

20

The defining name of a function that is a compilation unit shall not be an `operator_symbol`.

21

*Static Semantics*

A `subprogram_renaming_declaration` that is a `library_unit_renaming_declaration` is a renaming-as-declaration, not a renaming-as-body.

22

There are two kinds of dependences among compilation units:

23

- The *semantic dependences* (see below) are the ones needed to check the compile-time rules across compilation unit boundaries; a compilation unit depends semantically on the other compilation units needed to determine its legality. The visibility rules are based on the semantic dependences.
- The *elaboration dependences* (see 10.2) determine the order of elaboration of `library_items`.

24

25

A `library_item` depends semantically upon its parent declaration. A subunit depends semantically upon its parent body. A `library_unit_body` depends semantically upon the corresponding `library_unit_declaration`, if any. The declaration of the limited view of a library package depends semantically upon the declaration of the limited view of its parent. The declaration of a library package depends semantically upon the declaration of its limited view. A compilation unit depends semantically upon each `library_item` mentioned in a `with_clause` of the compilation unit. In addition, if a given compilation unit contains an `attribute_reference` of a type defined in another compilation unit, then the given compilation unit depends semantically upon the other compilation unit. The semantic dependence relationship is transitive.

26/2

*Dynamic Semantics*

- 26.1/2 The elaboration of the declaration of the limited view of a package has no effect.

## NOTES

27 1 A simple program may consist of a single compilation unit. A compilation need not have any compilation units; for example, its text can consist of pragmas.

28 2 The designator of a library function cannot be an operator\_symbol, but a nonlibrary renaming\_declaration is allowed to rename a library function as an operator. Within a partition, two library subprograms are required to have distinct names and hence cannot overload each other. However, renaming\_declarations are allowed to define overloaded names for such subprograms, and a locally declared subprogram is allowed to overload a library subprogram. The expanded name Standard.L can be used to denote a root library unit L (unless the declaration of Standard is hidden) since root library unit declarations occur immediately within the declarative region of package Standard.

*Examples*

## 29 Examples of library units:

```
30 package Rational_Numbers.IO is -- public child of Rational_Numbers, see 7.1
 procedure Put(R : in Rational);
 procedure Get(R : out Rational);
end Rational_Numbers.IO;

31 private procedure Rational_Numbers.Reduce(R : in out Rational); -- private child of Rational_Numbers

32 with Rational_Numbers.Reduce; -- refer to a private child
package body Rational_Numbers is
 ...
end Rational_Numbers;

33 with Rational_Numbers.IO; use Rational_Numbers;
with Ada.Text_Io; -- see A.7.0
procedure Main is -- a root library procedure
 R : Rational;
begin
 R := 5/3; -- construct a rational number, see 7.1
 Ada.Text_Io.Put("The answer is: ");
 IO.Put(R);
 Ada.Text_Io.New_Line;
end Main;

34 with Rational_Numbers.IO;
package Rational_IO renames Rational_Numbers.IO; -- a library unit renaming declaration
```

35 Each of the above library\_items can be submitted to the compiler separately.

## 10.1.2 Context Clauses - With Clauses

- 1 A context\_clause is used to specify the library\_items whose names are needed within a compilation unit.

*Syntax*

2 context\_clause ::= {context\_item}

3 context\_item ::= with\_clause | use\_clause

4/2 with\_clause ::= limited\_with\_clause | nonlimited\_with\_clause

4.1/2 limited\_with\_clause ::= limited [private] with library\_unit\_name {, library\_unit\_name};

4.2/2 nonlimited\_with\_clause ::= [private] with library\_unit\_name {, library\_unit\_name};

*Name Resolution Rules*

The *scope* of a *with\_clause* that appears on a *library\_unit\_declaration* or *library\_unit\_renaming\_declaration* consists of the entire declarative region of the declaration, which includes all children and subunits. The scope of a *with\_clause* that appears on a body consists of the body, which includes all subunits.

A *library\_item* (and the corresponding library unit) is *named* in a *with\_clause* if it is denoted by a *library\_unit\_name* in the *with\_clause*. A *library\_item* (and the corresponding library unit) is *mentioned* in a *with\_clause* if it is named in the *with\_clause* or if it is denoted by a prefix in the *with\_clause*.

Outside its own declarative region, the declaration or renaming of a library unit can be visible only within the scope of a *with\_clause* that mentions it. The visibility of the declaration or renaming of a library unit otherwise follows from its placement in the environment.

*Legality Rules*

If a *with\_clause* of a given *compilation\_unit* mentions a private child of some library unit, then the given *compilation\_unit* shall be one of:

- the declaration, body, or subunit of a private descendant of that library unit;
- the body or subunit of a public descendant of that library unit, but not a subprogram body acting as a subprogram declaration (see 10.1.4); or
- the declaration of a public descendant of that library unit, in which case the *with\_clause* shall include the reserved word **private**.

A name denoting a library item that is visible only due to being mentioned in one or more *with\_clauses* that include the reserved word **private** shall appear only within:

- a private part;
- a body, but not within the *subprogram\_specification* of a library subprogram body;
- a private descendant of the unit on which one of these *with\_clauses* appear; or
- a pragma within a context clause.

A *library\_item* mentioned in a *limited\_with\_clause* shall be the implicit declaration of the limited view of a library package, not the declaration of a subprogram, generic unit, generic instance, or a renaming.

A *limited\_with\_clause* shall not appear on a *library\_unit\_body*, *subunit*, or *library\_unit\_renaming\_declaration*.

A *limited\_with\_clause* that names a library package shall not appear:

- in the *context\_clause* for the explicit declaration of the named library package;
- in the same *context\_clause* as, or within the scope of, a *nonlimited\_with\_clause* that mentions the same library package; or
- in the same *context\_clause* as, or within the scope of, a *use\_clause* that names an entity declared within the declarative region of the library package.

## NOTES

3 A *library\_item* mentioned in a *nonlimited\_with\_clause* of a compilation unit is visible within the compilation unit and hence acts just like an ordinary declaration. Thus, within a compilation unit that mentions its declaration, the name of a library package can be given in *use\_clauses* and can be used to form expanded names, a library subprogram can be called, and instances of a generic library unit can be declared. If a child of a parent generic package is mentioned in a *nonlimited\_with\_clause*, then the corresponding declaration nested within each visible instance is visible within the

compilation unit. Similarly, a `library_item` mentioned in a `limited_with_clause` of a compilation unit is visible within the compilation unit and thus can be used to form expanded names.

#### Examples

```

24/2 package Office is
 end Office;

25/2 with Ada.Strings.Unbounded;
 package Office.Locations is
 type Location is new Ada.Strings.Unbounded.Unbounded_String;
 end Office.Locations;

26/2 limited with Office.Departments; -- types are incomplete
 private with Office.Locations; -- only visible in private part
 package Office.Employees is
 type Employee is private;
 function Dept_Of(Emp : Employee) return access Departments.Department;
 procedure Assign_Dept(Emp : in out Employee;
 Dept : access Departments.Department);
 ...
28/2 private
 type Employee is
 record
 Dept : access Departments.Department;
 Loc : Locations.Location;
 ...
 end record;
 end Office.Employees;

29/2 limited with Office.Employees;
 package Office.Departments is
 type Department is private;
 function Manager_Of(Dept : Department) return access Employees.Employee;
 procedure Assign_Manager(Dept : in out Department;
 Mgr : access Employees.Employee);
 ...
 end Office.Departments;

```

- 31/2 The `limited_with_clause` may be used to support mutually dependent abstractions that are split across multiple packages. In this case, an employee is assigned to a department, and a department has a manager who is an employee. If a `with_clause` with the reserved word `private` appears on one library unit and mentions a second library unit, it provides visibility to the second library unit, but restricts that visibility to the private part and body of the first unit. The compiler checks that no use is made of the second unit in the visible part of the first unit.

### 10.1.3 Subunits of Compilation Units

- 1 Subunits are like child units, with these (important) differences: subunits support the separate compilation of bodies only (not declarations); the parent contains a `body_stub` to indicate the existence and place of each of its subunits; declarations appearing in the parent's body can be visible within the subunits.

#### Syntax

```

2 body_stub ::= subprogram_body_stub | package_body_stub | task_body_stub | protected_body_stub
3/2 subprogram_body_stub ::= [overriding_indicator] subprogram_specification is separate;
4 package_body_stub ::= package body defining_identifier is separate;

```

|                                                                                        |   |
|----------------------------------------------------------------------------------------|---|
| task_body_stub ::= <b>task body</b> defining_identifier <b>is separate</b> ;           | 5 |
| protected_body_stub ::= <b>protected body</b> defining_identifier <b>is separate</b> ; | 6 |
| subunit ::= <b>separate</b> (parent_unit_name) proper_body                             | 7 |

*Legality Rules*

The *parent body* of a subunit is the body of the program unit denoted by its *parent\_unit\_name*. The term *subunit* is used to refer to a subunit and also to the *proper\_body* of a subunit. The *subunits of a program unit* include any subunit that names that program unit as its parent, as well as any subunit that names such a subunit as its parent (recursively).

The parent body of a subunit shall be present in the current environment, and shall contain a corresponding *body\_stub* with the same *defining\_identifier* as the subunit.

A *package\_body\_stub* shall be the completion of a *package\_declaration* or *generic\_package\_declaration*; a *task\_body\_stub* shall be the completion of a task declaration; a *protected\_body\_stub* shall be the completion of a protected declaration.

In contrast, a *subprogram\_body\_stub* need not be the completion of a previous declaration, in which case the *\_stub* declares the subprogram. If the *\_stub* is a completion, it shall be the completion of a *subprogram\_declaration* or *generic\_subprogram\_declaration*. The profile of a *subprogram\_body\_stub* that completes a declaration shall conform fully to that of the declaration.

A subunit that corresponds to a *body\_stub* shall be of the same kind (*package\_*, *subprogram\_*, *task\_*, or *protected\_*) as the *body\_stub*. The profile of a *subprogram\_body* subunit shall be fully conformant to that of the corresponding *body\_stub*.

A *body\_stub* shall appear immediately within the *declarative\_part* of a compilation unit body. This rule does not apply within an instance of a generic unit.

The *defining\_identifiers* of all *body\_stubs* that appear immediately within a particular *declarative\_part* shall be distinct.

*Post-Compilation Rules*

For each *body\_stub*, there shall be a subunit containing the corresponding *proper\_body*.

NOTES

4 The rules in 10.1.4, “The Compilation Process” say that a *body\_stub* is equivalent to the corresponding *proper\_body*. This implies:

- Visibility within a subunit is the visibility that would be obtained at the place of the corresponding *body\_stub* (within the parent body) if the *context\_clause* of the subunit were appended to that of the parent body.
- The effect of the elaboration of a *body\_stub* is to elaborate the subunit.

*Examples*

The package Parent is first written without subunits:

```
package Parent is
 procedure Inner;
end Parent;

with Ada.Text_IO;
package body Parent is
 Variable : String := "Hello, there.";
 procedure Inner is
 begin
 Ada.Text_IO.Put_Line(Variable);
 end Inner;
end Parent;
```

- 22 The body of procedure Inner may be turned into a subunit by rewriting the package body as follows (with the declaration of Parent remaining the same):

```

23 package body Parent is
24 Variable : String := "Hello, there.";
25 procedure Inner is separate;
26 end Parent;
27
28 with Ada.Text_IO;
29 separate(Parent)
30 procedure Inner is
31 begin
32 Ada.Text_IO.Put_Line(Variable);
33 end Inner;

```

## 10.1.4 The Compilation Process

- 1 Each compilation unit submitted to the compiler is compiled in the context of an *environment\_declarative\_part* (or simply, an *environment*), which is a conceptual *declarative\_part* that forms the outermost declarative region of the context of any *compilation*. At run time, an environment forms the *declarative\_part* of the body of the environment task of a partition (see 10.2, “Program Execution”).
- 2 The *declarative\_items* of the environment are *library\_items* appearing in an order such that there are no forward semantic dependences. Each included subunit occurs in place of the corresponding stub. The visibility rules apply as if the environment were the outermost declarative region, except that *with\_clauses* are needed to make declarations of library units visible (see 10.1.2).
- 3/2 The mechanisms for creating an environment and for adding and replacing compilation units within an environment are implementation defined. The mechanisms for adding a compilation unit mentioned in a *limited\_with\_clause* to an environment are implementation defined.

### *Name Resolution Rules*

- 4/1 If a *library\_unit\_body* that is a *subprogram\_body* is submitted to the compiler, it is interpreted only as a completion if a *library\_unit\_declaration* with the same *defining\_program\_unit\_name* already exists in the environment for a subprogram other than an instance of a generic subprogram or for a generic subprogram (even if the profile of the body is not type conformant with that of the declaration); otherwise the *subprogram\_body* is interpreted as both the declaration and body of a library subprogram.

### *Legality Rules*

- 5 When a compilation unit is compiled, all compilation units upon which it depends semantically shall already exist in the environment; the set of these compilation units shall be *consistent* in the sense that the new compilation unit shall not semantically depend (directly or indirectly) on two different versions of the same compilation unit, nor on an earlier version of itself.

### *Implementation Permissions*

- 6/2 The implementation may require that a compilation unit be legal before it can be mentioned in a *limited\_with\_clause* or it can be inserted into the environment.
- 7/2 When a compilation unit that declares or renames a library unit is added to the environment, the implementation may remove from the environment any preexisting *library\_item* or subunit with the same full expanded name. When a compilation unit that is a subunit or the body of a library unit is added to the environment, the implementation may remove from the environment any preexisting version of the same compilation unit. When a compilation unit that contains a *body\_stub* is added to the environment, the implementation may remove any preexisting *library\_item* or subunit with the same full expanded name as

the body\_stub. When a given compilation unit is removed from the environment, the implementation may also remove any compilation unit that depends semantically upon the given one. If the given compilation unit contains the body of a subprogram to which a **pragma** **Inline** applies, the implementation may also remove any compilation unit containing a call to that subprogram.

## NOTES

5 The rules of the language are enforced across compilation and compilation unit boundaries, just as they are enforced within a single compilation unit. 8

6 An implementation may support a concept of a *library*, which contains *library\_items*. If multiple libraries are supported, the implementation has to define how a single environment is constructed when a compilation unit is submitted to the compiler. Naming conflicts between different libraries might be resolved by treating each library as the root of a hierarchy of child library units. 9

7 A compilation unit containing an instantiation of a separately compiled generic unit does not semantically depend on the body of the generic unit. Therefore, replacing the generic body in the environment does not result in the removal of the compilation unit containing the instantiation. 10

## 10.1.5 Pragmas and Program Units

This subclause discusses pragmas related to program units, library units, and compilations. 1

### *Name Resolution Rules*

Certain **pragmas** are defined to be *program unit pragmas*. A **name** given as the argument of a program unit **pragma** shall resolve to denote the declarations or renamings of one or more program units that occur immediately within the declarative region or compilation in which the **pragma** immediately occurs, or it shall resolve to denote the declaration of the immediately enclosing program unit (if any); the **pragma** applies to the denoted program unit(s). If there are no **names** given as arguments, the **pragma** applies to the immediately enclosing program unit. 2

### *Legality Rules*

A program unit **pragma** shall appear in one of these places: 3

- At the place of a *compilation\_unit*, in which case the **pragma** shall immediately follow in the same compilation (except for other **pragmas**) a *library\_unit\_declaration* that is a *subprogram\_declaration*, *generic\_subprogram\_declaration*, or *generic\_instantiation*, and the **pragma** shall have an argument that is a **name** denoting that declaration. 4
- Immediately within the visible part of a program unit and before any nested declaration (but not within a generic formal part), in which case the argument, if any, shall be a *direct\_name* that denotes the immediately enclosing program unit declaration. 5/1
- At the place of a declaration other than the first, of a *declarative\_part* or program unit declaration, in which case the **pragma** shall have an argument, which shall be a *direct\_name* that denotes one or more of the following (and nothing else): a *subprogram\_declaration*, a *generic\_subprogram\_declaration*, or a *generic\_instantiation*, of the same *declarative\_part* or program unit declaration. 6

Certain program unit **pragmas** are defined to be *library unit pragmas*. The **name**, if any, in a library unit **pragma** shall denote the declaration of a library unit. 7

### *Static Semantics*

A library unit **pragma** that applies to a generic unit does not apply to its instances, unless a specific rule for the **pragma** specifies the contrary. 7.1/1

*Post-Compilation Rules*

- 8 Certain pragmas are defined to be *configuration pragmas*; they shall appear before the first compilation\_unit of a compilation. They are generally used to select a partition-wide or system-wide option. The pragma applies to all compilation\_units appearing in the compilation, unless there are none, in which case it applies to all future compilation\_units compiled into the same environment.

*Implementation Permissions*

- 9/2 An implementation may require that configuration pragmas that select partition-wide or system-wide options be compiled when the environment contains no library\_items other than those of the predefined environment. In this case, the implementation shall still accept configuration pragmas in individual compilations that confirm the initially selected partition-wide or system-wide options.

*Implementation Advice*

- 10/1 When applied to a generic unit, a program unit pragma that is not a library unit pragma should apply to each instance of the generic unit for which there is not an overriding pragma applied directly to the instance.

## 10.1.6 Environment-Level Visibility Rules

- 1 The normal visibility rules do not apply within a parent\_unit\_name or a context\_clause, nor within a pragma that appears at the place of a compilation unit. The special visibility rules for those contexts are given here.

*Static Semantics*

- 2/2 Within the parent\_unit\_name at the beginning of an explicit library\_item, and within a nonlimited\_with\_clause, the only declarations that are visible are those that are explicit library\_items of the environment, and the only declarations that are directly visible are those that are explicit root library\_items of the environment. Within a limited\_with\_clause, the only declarations that are visible are those that are the implicit declaration of the limited view of a library package of the environment, and the only declarations that are directly visible are those that are the implicit declaration of the limited view of a root library package.

- 3 Within a use\_clause or pragma that is within a context\_clause, each library\_item mentioned in a previous with\_clause of the same context\_clause is visible, and each root library\_item so mentioned is directly visible. In addition, within such a use\_clause, if a given declaration is visible or directly visible, each declaration that occurs immediately within the given declaration's visible part is also visible. No other declarations are visible or directly visible.

- 4 Within the parent\_unit\_name of a subunit, library\_items are visible as they are in the parent\_unit\_name of a library\_item; in addition, the declaration corresponding to each body\_stub in the environment is also visible.

- 5 Within a pragma that appears at the place of a compilation unit, the immediately preceding library\_item and each of its ancestors is visible. The ancestor root library\_item is directly visible.

- 6/2 Notwithstanding the rules of 4.1.3, an expanded name in a with\_clause, a pragma in a context\_clause, or a pragma that appears at the place of a compilation unit may consist of a prefix that denotes a generic package and a selector\_name that denotes a child of that generic package. (The child is necessarily a generic unit; see 10.1.1.)

## 10.2 Program Execution

An Ada *program* consists of a set of *partitions*, which can execute in parallel with one another, possibly in a separate address space, and possibly on a separate computer.

### Post-Compilation Rules

A partition is a program or part of a program that can be invoked from outside the Ada implementation. For example, on many systems, a partition might be an executable file generated by the system linker. The user can *explicitly assign* library units to a partition. The assignment is done in an implementation-defined manner. The compilation units included in a partition are those of the explicitly assigned library units, as well as other compilation units *needed* by those library units. The compilation units needed by a given compilation unit are determined as follows (unless specified otherwise via an implementation-defined pragma, or by some other implementation-defined means):

- A compilation unit needs itself; 3
- If a compilation unit is needed, then so are any compilation units upon which it depends semantically; 4
- If a `library_unit_declaration` is needed, then so is any corresponding `library_unit_body`; 5
- If a compilation unit with stubs is needed, then so are any corresponding subunits; 6/2
- If the (implicit) declaration of the limited view of a library package is needed, then so is the explicit declaration of the library package. 6.1/2

The user can optionally designate (in an implementation-defined manner) one subprogram as the *main subprogram* for the partition. A main subprogram, if specified, shall be a subprogram. 7

Each partition has an anonymous *environment task*, which is an implicit outermost task whose execution elaborates the `library_items` of the environment `declarative_part`, and then calls the main subprogram, if there is one. A partition's execution is that of its tasks. 8

The order of elaboration of library units is determined primarily by the *elaboration dependences*. There is an elaboration dependence of a given `library_item` upon another if the given `library_item` or any of its subunits depends semantically on the other `library_item`. In addition, if a given `library_item` or any of its subunits has a pragma `Elaborate` or `Elaborate_All` that names another library unit, then there is an elaboration dependence of the given `library_item` upon the body of the other library unit, and, for `Elaborate_All` only, upon each `library_item` needed by the declaration of the other library unit. 9

The environment task for a partition has the following structure:

```
task Environment_Task;
task body Environment_Task is
 ... (1) -- The environment declarative_part
 -- (that is, the sequence of library_items) goes here.
begin
 ... (2) -- Call the main subprogram, if there is one.
end Environment_Task;
```

The environment `declarative_part` at (1) is a sequence of `declarative_items` consisting of copies of the `library_items` included in the partition. The order of elaboration of `library_items` is the order in which they appear in the environment `declarative_part`: 13

- The order of all included `library_items` is such that there are no forward elaboration dependences. 14

- 15     • Any included `library_unit_declaration` to which a `pragma Elaborate_Body` applies is  
      immediately followed by its `library_unit_body`, if included.
- 16     • All `library_items` declared pure occur before any that are not declared pure.
- 17     • All preelaborated `library_items` occur before any that are not preelaborated.
- 18     There shall be a total order of the `library_items` that obeys the above rules. The order is otherwise  
      implementation defined.
- 19     The full expanded names of the library units and subunits included in a given partition shall be distinct.
- 20     The `sequence_of_statements` of the environment task (see (2) above) consists of either:
- 21         • A call to the main subprogram, if the partition has one. If the main subprogram has parameters,  
          they are passed; where the actuals come from is implementation defined. What happens to the  
          result of a main function is also implementation defined.
- 22         or:
- 23         • A `null_statement`, if there is no main subprogram.
- 24     The mechanisms for building and running partitions are implementation defined. These might be  
      combined into one operation, as, for example, in dynamic linking, or “load-and-go” systems.

*Dynamic Semantics*

- 25     The execution of a program consists of the execution of a set of partitions. Further details are  
      implementation defined. The execution of a partition starts with the execution of its environment task,  
      ends when the environment task terminates, and includes the executions of all tasks of the partition. The  
      execution of the (implicit) `task_body` of the environment task acts as a master for all other tasks created as  
      part of the execution of the partition. When the environment task completes (normally or abnormally), it  
      waits for the termination of all such tasks, and then finalizes any remaining objects of the partition.

*Bounded (Run-Time) Errors*

- 26     Once the environment task has awaited the termination of all other tasks of the partition, any further  
      attempt to create a task (during finalization) is a bounded error, and may result in the raising of  
      `Program_Error` either upon creation or activation of the task. If such a task is activated, it is not specified  
      whether the task is awaited prior to termination of the environment task.

*Implementation Requirements*

- 27     The implementation shall ensure that all compilation units included in a partition are consistent with one  
      another, and are legal according to the rules of the language.

*Implementation Permissions*

- 28     The kind of partition described in this clause is known as an *active* partition. An implementation is  
      allowed to support other kinds of partitions, with implementation-defined semantics.
- 29     An implementation may restrict the kinds of subprograms it supports as main subprograms. However, an  
      implementation is required to support all main subprograms that are public parameterless library  
      procedures.
- 30     If the environment task completes abnormally, the implementation may abort any dependent tasks.

**NOTES**

- 31     8 An implementation may provide inter-partition communication mechanism(s) via special packages and pragmas.  
      Standard pragmas for distribution and methods for specifying inter-partition communication are defined in Annex E,

“Distributed Systems”. If no such mechanisms are provided, then each partition is isolated from all others, and behaves as a program in and of itself.

9 Partitions are not required to run in separate address spaces. For example, an implementation might support dynamic linking via the partition concept. 32

10 An order of elaboration of *library\_items* that is consistent with the partial ordering defined above does not always ensure that each *library\_unit\_body* is elaborated before any other compilation unit whose elaboration necessitates that the *library\_unit\_body* be already elaborated. (In particular, there is no requirement that the body of a library unit be elaborated as soon as possible after the *library\_unit\_declaration* is elaborated, unless the pragmas in subclause 10.2.1 are used.) 33

11 A partition (active or otherwise) need not have a main subprogram. In such a case, all the work done by the partition would be done by elaboration of various *library\_items*, and by tasks created by that elaboration. Passive partitions, which cannot have main subprograms, are defined in Annex E, “Distributed Systems”. 34

## 10.2.1 Elaboration Control

This subclause defines pragmas that help control the elaboration order of *library\_items*. 1

### Syntax

The form of a **pragma** Preelaborate is as follows: 2

**pragma** Preelaborate[*(library\_unit\_name)*]; 3

A **pragma** Preelaborate is a library unit pragma. 4

The form of a **pragma** Preelaborable\_Initialization is as follows: 4.1/2

**pragma** Preelaborable\_Initialization(*direct\_name*); 4.2/2

### Legality Rules

An elaborable construct is preelaborable unless its elaboration performs any of the following actions: 5

- The execution of a **statement** other than a **null\_statement**. 6
- A call to a subprogram other than a static function. 7
- The evaluation of a primary that is a **name** of an object, unless the **name** is a static expression, or statically denotes a discriminant of an enclosing type. 8
- The creation of an object (including a component) of a type that does not have preelaborable initialization. Similarly, the evaluation of an **extension\_aggregate** with an ancestor **subtype\_mark** denoting a subtype of such a type. 9/2

A generic body is preelaborable only if elaboration of a corresponding instance body would not perform any such actions, presuming that: 10/2

- the actual for each formal private type (or extension) declared within the formal part of the generic unit is a private type (or extension) that does not have preelaborable initialization; 10.1/2
- the actual for each formal type is nonstatic; 10.2/2
- the actual for each formal object is nonstatic; and 10.3/2
- the actual for each formal subprogram is a user-defined subprogram. 10.4/2

If a **pragma** Preelaborate (or **pragma** Pure — see below) applies to a library unit, then it is *preelaborated*. 11/1  
If a library unit is preelaborated, then its declaration, if any, and body, if any, are elaborated prior to all non-preelaborated *library\_items* of the partition. The declaration and body of a preelaborated library unit, and all subunits that are elaborated as part of elaborating the library unit, shall be preelaborable. In addition to the places where Legality Rules normally apply (see 12.3), this rule applies also in the private

part of an instance of a generic unit. In addition, all compilation units of a preelaborated library unit shall depend semantically only on compilation units of other preelaborated library units.

- 11.1/2 The following rules specify which entities have *preelaborable initialization*:
- 11.2/2
  - The partial view of a private type or private extension, a protected type without `entry_declarations`, a generic formal private type, or a generic formal derived type, have preelaborable initialization if and only if the pragma `Preeelaborable_Initialization` has been applied to them. A protected type with `entry_declarations` or a task type never has preelaborable initialization.
- 11.3/2
  - A component (including a discriminant) of a record or protected type has preelaborable initialization if its declaration includes a `default_expression` whose execution does not perform any actions prohibited in preelaborable constructs as described above, or if its declaration does not include a default expression and its type has preelaborable initialization.
- 11.4/2
  - A derived type has preelaborable initialization if its parent type has preelaborable initialization and (in the case of a derived record extension) if the non-inherited components all have preelaborable initialization. However, a user-defined controlled type with an overriding `Initialize` procedure does not have preelaborable initialization.
- 11.5/2
  - A view of a type has preelaborable initialization if it is an elementary type, an array type whose component type has preelaborable initialization, a record type whose components all have preelaborable initialization, or an interface type.
- 11.6/2 A pragma `Preeelaborable_Initialization` specifies that a type has preelaborable initialization. This pragma shall appear in the visible part of a package or generic package.
- 11.7/2 If the pragma appears in the first list of `basic_declarative_items` of a `package_specification`, then the `direct_name` shall denote the first subtype of a private type, private extension, or protected type that is not an interface type and is without `entry_declarations`, and the type shall be declared immediately within the same package as the pragma. If the pragma is applied to a private type or a private extension, the full view of the type shall have preelaborable initialization. If the pragma is applied to a protected type, each component of the protected type shall have preelaborable initialization. In addition to the places where Legality Rules normally apply, these rules apply also in the private part of an instance of a generic unit.
- 11.8/2 If the pragma appears in a `generic_formal_part`, then the `direct_name` shall denote a generic formal private type or a generic formal derived type declared in the same `generic_formal_part` as the pragma. In a `generic_instantiation` the corresponding actual type shall have preelaborable initialization.

#### *Implementation Advice*

- 12 In an implementation, a type declared in a preelaborated package should have the same representation in every elaboration of a given version of the package, whether the elaborations occur in distinct executions of the same program, or in executions of distinct programs or partitions that include the given version.

#### *Syntax*

- 13 The form of a `pragma Pure` is as follows:
- 14   `pragma Pure[(library_unit_name)];`
- 15 A `pragma Pure` is a library unit pragma.

#### *Static Semantics*

- 15.1/2 A *pure library\_item* is a preelaborable *library\_item* whose elaboration does not perform any of the following actions:
  - the elaboration of a variable declaration;

- the evaluation of an allocator of an access-to-variable type; for the purposes of this rule, the partial view of a type is presumed to have non-visible components whose default initialization evaluates such an allocator; 15.3/2
- the elaboration of the declaration of a named access-to-variable type unless the Storage\_Size of the type has been specified by a static expression with value zero or is defined by the language to be zero; 15.4/2
- the elaboration of the declaration of a named access-to-constant type for which the Storage\_Size has been specified by an expression other than a static expression with value zero. 15.5/2

The Storage\_Size for an anonymous access-to-variable type declared at library level in a library unit that is declared pure is defined to be zero. 15.6/2

#### *Legality Rules*

*This paragraph was deleted.*

16/2

A **pragma Pure** is used to declare that a library unit is pure. If a **pragma Pure** applies to a library unit, then its compilation units shall be pure, and they shall depend semantically only on compilation units of other library units that are declared pure. Furthermore, the full view of any partial view declared in the visible part of the library unit that has any available stream attributes shall support external streaming (see 13.13.2). 17/2

#### *Implementation Permissions*

If a library unit is declared pure, then the implementation is permitted to omit a call on a library-level subprogram of the library unit if the results are not needed after the call. In addition, the implementation may omit a call on such a subprogram and simply reuse the results produced by an earlier call on the same subprogram, provided that none of the parameters nor any object accessible via access values from the parameters are of a limited type, and the addresses and values of all by-reference actual parameters, the values of all by-copy-in actual parameters, and the values of all objects accessible via access values from the parameters, are the same as they were at the earlier call. This permission applies even if the subprogram produces other side effects when called. 18/2

#### *Syntax*

The form of a **pragma Elaborate**, **Elaborate\_All**, or **Elaborate\_Body** is as follows:

19

**pragma** Elaborate(*library\_unit\_name*{, *library\_unit\_name*});

20

**pragma** Elaborate\_All(*library\_unit\_name*{, *library\_unit\_name*});

21

**pragma** Elaborate\_Body[*(library\_unit\_name)*];

22

A **pragma Elaborate** or **Elaborate\_All** is only allowed within a **context\_clause**.

23

A **pragma Elaborate\_Body** is a library unit pragma.

24

#### *Legality Rules*

If a **pragma Elaborate\_Body** applies to a declaration, then the declaration requires a completion (a body).

25

The *library\_unit\_name* of a **pragma Elaborate** or **Elaborate\_All** shall denote a nonlimited view of a library unit. 25.1/2

#### *Static Semantics*

A **pragma Elaborate** specifies that the body of the named library unit is elaborated before the current *library\_item*. A **pragma Elaborate\_All** specifies that each *library\_item* that is needed by the named library

26

unit declaration is elaborated before the current `library_item`. A `pragma Elaborate_Body` specifies that the body of the library unit is elaborated immediately after its declaration.

NOTES

- 27    12 A preelaborated library unit is allowed to have non-preelaborable children.  
28    13 A library unit that is declared pure is allowed to have impure children.

## Section 11: Exceptions

This section defines the facilities for dealing with errors or other exceptional situations that arise during program execution. An *exception* represents a kind of exceptional situation; an occurrence of such a situation (at run time) is called an *exception occurrence*. To *raise* an exception is to abandon normal program execution so as to draw attention to the fact that the corresponding situation has arisen. Performing some actions in response to the arising of an exception is called *handling* the exception.

An *exception\_declaration* declares a name for an exception. An exception is raised initially either by a *raise\_statement* or by the failure of a language-defined check. When an exception arises, control can be transferred to a user-provided *exception\_handler* at the end of a *handled\_sequence\_of\_statements*, or it can be propagated to a dynamically enclosing execution.

### 11.1 Exception Declarations

An *exception\_declaration* declares a name for an exception.

#### Syntax

```
exception_declaration ::= defining_identifier_list : exception;
```

#### Static Semantics

Each single *exception\_declaration* declares a name for a different exception. If a generic unit includes an *exception\_declaration*, the *exception\_declarations* implicitly generated by different instantiations of the generic unit refer to distinct exceptions (but all have the same *defining\_identifier*). The particular exception denoted by an exception name is determined at compilation time and is the same regardless of how many times the *exception\_declaration* is elaborated.

The *predefined* exceptions are the ones declared in the declaration of package Standard: Constraint\_Error, Program\_Error, Storage\_Error, and Tasking\_Error; one of them is raised when a language-defined check fails.

#### Dynamic Semantics

The elaboration of an *exception\_declaration* has no effect.

The execution of any construct raises Storage\_Error if there is insufficient storage for that execution. The amount of storage needed for the execution of constructs is unspecified.

#### Examples

*Examples of user-defined exception declarations:*

```
Singular : exception;
Error : exception;
Overflow, Underflow : exception;
```

### 11.2 Exception Handlers

The response to one or more exceptions is specified by an *exception\_handler*.

*Syntax*

```

2 handled_sequence_of_statements ::= sequence_of_statements
3 [exception exception_handler {exception_handler}]
4 exception_handler ::= when [choice_parameter_specification:] exception_choice { exception_choice } => sequence_of_statements
5 choice_parameter_specification ::= defining_identifier
6 exception_choice ::= exception_name | others

```

*Legality Rules*

- 6 A choice with an *exception\_name* covers the named exception. A choice with **others** covers all exceptions not named by previous choices of the same *handled\_sequence\_of\_statements*. Two choices in different *exception\_handlers* of the same *handled\_sequence\_of\_statements* shall not cover the same exception.
- 7 A choice with **others** is allowed only for the last handler of a *handled\_sequence\_of\_statements* and as the only choice of that handler.
- 8 An *exception\_name* of a choice shall not denote an exception declared in a generic formal package.

*Static Semantics*

- 9 A *choice\_parameter\_specification* declares a *choice parameter*, which is a constant object of type *Exception\_Occurrence* (see 11.4.1). During the handling of an exception occurrence, the choice parameter, if any, of the handler represents the exception occurrence that is being handled.

*Dynamic Semantics*

- 10 The execution of a *handled\_sequence\_of\_statements* consists of the execution of the *sequence\_of\_statements*. The optional handlers are used to handle any exceptions that are propagated by the *sequence\_of\_statements*.

*Examples*

```

11 Example of an exception handler:
12 begin
 Open(File, In_File, "input.txt"); -- see A.8.2
 exception
 when E : Name_Error =>
 Put("Cannot open input file : ");
 Put_Line(Exception_Message(E)); -- see 11.4.1
 raise;
 end;

```

## 11.3 Raise Statements

- 1 A *raise\_statement* raises an exception.

*Syntax*

```

2/2 raise_statement ::= raise;
 | raise exception_name [with string_expression];

```

*Legality Rules*

The name, if any, in a `raise_statement` shall denote an exception. A `raise_statement` with no `exception_name` (that is, a *re-raise statement*) shall be within a handler, but not within a body enclosed by that handler.

*Name Resolution Rules*

The expression, if any, in a `raise_statement`, is expected to be of type String.

3

3.1/2

*Dynamic Semantics*

To *raise an exception* is to raise a new occurrence of that exception, as explained in 11.4. For the execution of a `raise_statement` with an `exception_name`, the named exception is raised. If a `string_expression` is present, the `expression` is evaluated and its value is associated with the exception occurrence. For the execution of a re-raise statement, the exception occurrence that caused transfer of control to the innermost enclosing handler is raised again.

4/2

*Examples*

*Examples of raise statements:*

5

```
raise Ada.IO_Exceptions.Name_Error; -- see A.13
raise Queue_Error with "Buffer Full"; -- see 9.11
raise; -- re-raise the current exception
```

6/2

7

## 11.4 Exception Handling

When an exception occurrence is raised, normal program execution is abandoned and control is transferred to an applicable `exception_handler`, if any. To *handle* an exception occurrence is to respond to the exceptional event. To *propagate* an exception occurrence is to raise it again in another context; that is, to fail to respond to the exceptional event in the present context.

1

*Dynamic Semantics*

Within a given task, if the execution of construct *a* is defined by this International Standard to consist (in part) of the execution of construct *b*, then while *b* is executing, the execution of *a* is said to *dynamically enclose* the execution of *b*. The *innermost dynamically enclosing* execution of a given execution is the dynamically enclosing execution that started most recently.

2

When an exception occurrence is raised by the execution of a given construct, the rest of the execution of that construct is *abandoned*; that is, any portions of the execution that have not yet taken place are not performed. The construct is first completed, and then left, as explained in 7.6.1. Then:

3

- If the construct is a `task_body`, the exception does not propagate further;
- If the construct is the `sequence_of_statements` of a `handled_sequence_of_statements` that has a handler with a choice covering the exception, the occurrence is handled by that handler;
- Otherwise, the occurrence is *propagated* to the innermost dynamically enclosing execution, which means that the occurrence is raised again in that context.

4

5

6

When an occurrence is *handled* by a given handler, the `choice_parameter_specification`, if any, is first elaborated, which creates the choice parameter and initializes it to the occurrence. Then, the `sequence_of_statements` of the handler is executed; this execution replaces the abandoned portion of the execution of the `sequence_of_statements`.

7

## NOTES

- 8 1 Note that exceptions raised in a declarative\_part of a body are not handled by the handlers of the handled\_sequence\_of\_statements of that body.

## 11.4.1 The Package Exceptions

### *Static Semantics*

- 1 The following language-defined library package exists:

```

2/2 with Ada.Streams;
 package Ada.Exceptions is
 pragma Preelaborate(Exceptions);
 type Exception_Id is private;
 pragma Preelaborate_Initialization(Exception_Id);
 Null_Id : constant Exception_Id;
 function Exception_Name(Id : Exception_Id) return String;
 function Wide_Exception_Name(Id : Exception_Id) return Wide_String;
 function Wide_Wide_Exception_Name(Id : Exception_Id)
 return Wide_Wide_String;

3/2 type Exception_Occurrence is limited private;
 pragma Preelaborate_Initialization(Exception_Occurrence);
 type Exception_Occurrence_Access is access all Exception_Occurrence;
 Null_Occurrence : constant Exception_Occurrence;

4/2 procedure Raise_Exception(E : in Exception_Id;
 Message : in String := "");
 pragma No_Return(Raise_Exception);
 function Exception_Message(X : Exception_Occurrence) return String;
 procedure Reraise_Occurrence(X : in Exception_Occurrence);

5/2 function Exception_Identity(X : Exception_Occurrence)
 return Exception_Id;
 function Exception_Name(X : Exception_Occurrence) return String;
 -- Same as Exception_Name(Exception_Identity(X));
 function Wide_Exception_Name(X : Exception_Occurrence)
 return Wide_String;
 -- Same as Wide_Exception_Name(Exception_Identity(X));
 function Wide_Wide_Exception_Name(X : Exception_Occurrence)
 return Wide_Wide_String;
 -- Same as Wide_Wide_Exception_Name(Exception_Identity(X));
 function Exception_Information(X : Exception_Occurrence) return String;

6/2 procedure Save_Occurrence(Target : out Exception_Occurrence;
 Source : in Exception_Occurrence);
 function Save_Occurrence(Source : Exception_Occurrence)
 return Exception_Occurrence_Access;

6.1/2 procedure Read_Exception_Occurrence
 (Stream : not null access Ada.Streams.Root_Stream_Type'Class;
 Item : out Exception_Occurrence);
 procedure Write_Exception_Occurrence
 (Stream : not null access Ada.Streams.Root_Stream_Type'Class;
 Item : in Exception_Occurrence);

6.2/2 for Exception_Occurrence'Read use Read_Exception_Occurrence;
 for Exception_Occurrence'Write use Write_Exception_Occurrence;

6.3/2 private
 ... -- not specified by the language
 end Ada.Exceptions;

```

- 7 Each distinct exception is represented by a distinct value of type `Exception_Id`. `Null_Id` does not represent any exception, and is the default initial value of type `Exception_Id`. Each occurrence of an exception is represented by a value of type `Exception_Occurrence`. `Null_Occurrence` does not represent any exception occurrence, and is the default initial value of type `Exception_Occurrence`.

|                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                       |        |
|-----------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|--------|
| For a prefix E that denotes an exception, the following attribute is defined:                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                         | 8/1    |
| E'Identity      E'Identity returns the unique identity of the exception. The type of this attribute is Exception_Id.                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                  | 9      |
| Raise_Exception raises a new occurrence of the identified exception.                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                  | 10/2   |
| Exception_Message returns the message associated with the given Exception_Occurrence. For an occurrence raised by a call to Raise_Exception, the message is the Message parameter passed to Raise_Exception. For the occurrence raised by a raise_statement with an <i>exception_name</i> and a <i>string_expression</i> , the message is the <i>string_expression</i> . For the occurrence raised by a raise_statement with an <i>exception_name</i> but without a <i>string_expression</i> , the message is a string giving implementation-defined information about the exception occurrence. In all cases, Exception_Message returns a string with lower bound 1. | 10.1/2 |
| Reraise_Occurrence reraises the specified exception occurrence.                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                       | 10.2/2 |
| Exception_Identity returns the identity of the exception of the occurrence.                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                           | 11     |
| The Wide_Wide_Exception_Name functions return the full expanded name of the exception, in upper case, starting with a root library unit. For an exception declared immediately within package Standard, the defining_identifier is returned. The result is implementation defined if the exception is declared within an unnamed block_statement.                                                                                                                                                                                                                                                                                                                     | 12/2   |
| The Exception_Name functions (respectively, Wide_Exception_Name) return the same sequence of graphic characters as that defined for Wide_Wide_Exception_Name, if all the graphic characters are defined in Character (respectively, Wide_Character); otherwise, the sequence of characters is implementation defined, but no shorter than that returned by Wide_Wide_Exception_Name for the same value of the argument.                                                                                                                                                                                                                                               | 12.1/2 |
| The string returned by the Exception_Name, Wide_Exception_Name, and Wide_Wide_Exception_Name functions has lower bound 1.                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                             | 12.2/2 |
| Exception_Information returns implementation-defined information about the exception occurrence. The returned string has lower bound 1.                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                               | 13/2   |
| Reraise_Occurrence has no effect in the case of Null_Occurrence. Raise_Exception and Exception_Name raise Constraint_Error for a Null_Id. Exception_Message, Exception_Name, and Exception_Information raise Constraint_Error for a Null_Occurrence. Exception_Identity applied to Null_Occurrence returns Null_Id.                                                                                                                                                                                                                                                                                                                                                   | 14/2   |
| The Save_Occurrence procedure copies the Source to the Target. The Save_Occurrence function uses an allocator of type Exception_Occurrence_Access to create a new object, copies the Source to this new object, and returns an access value designating this new object; the result may be deallocated using an instance of Unchecked_Deallocation.                                                                                                                                                                                                                                                                                                                   | 15     |
| Write_Exception_Occurrence writes a representation of an exception occurrence to a stream; Read_Exception_Occurrence reconstructs an exception occurrence from a stream (including one written in a different partition).                                                                                                                                                                                                                                                                                                                                                                                                                                             | 15.1/2 |

*Implementation Requirements**This paragraph was deleted.*

16/2

*Implementation Permissions*

- 17 An implementation of `Exception_Name` in a space-constrained environment may return the `defining_-identifier` instead of the full expanded name.
- 18 The string returned by `Exception_Message` may be truncated (to no less than 200 characters) by the `Save_Occurrence` procedure (not the function), the `Reraise_Occurrence` procedure, and the re-raise statement.

*Implementation Advice*

- 19 `Exception_Message` (by default) and `Exception_Information` should produce information useful for debugging. `Exception_Message` should be short (about one line), whereas `Exception_Information` can be long. `Exception_Message` should not include the `Exception_Name`. `Exception_Information` should include both the `Exception_Name` and the `Exception_Message`.

## 11.4.2 Pragmas Assert and Assertion\_Policy

- 1/2 `Pragma Assert` is used to assert the truth of a Boolean expression at any point within a sequence of declarations or statements. `Pragma Assertion_Policy` is used to control whether such assertions are to be ignored by the implementation, checked at run-time, or handled in some implementation-defined manner.

*Syntax*

- 2/2 The form of a `pragma Assert` is as follows:
- ```
pragma Assert([Check =>] boolean_expression[, [Message =>] string_expression]);
```
- 4/2 A `pragma Assert` is allowed at the place where a `declarative_item` or a `statement` is allowed.
- 5/2 The form of a `pragma Assertion_Policy` is as follows:
- ```
pragma Assertion_Policy(policy_identifier);
```
- 7/2 A `pragma Assertion_Policy` is a configuration pragma.

*Name Resolution Rules*

- 8/2 The expected type for the `boolean_expression` of a `pragma Assert` is any boolean type. The expected type for the `string_expression` of a `pragma Assert` is type `String`.

*Legality Rules*

- 9/2 The `policy_identifier` of a `pragma Assertion_Policy` shall be either `Check`, `Ignore`, or an implementation-defined identifier.

*Static Semantics*

- 10/2 A `pragma Assertion_Policy` is a configuration pragma that specifies the assertion policy in effect for the compilation units to which it applies. Different policies may apply to different compilation units within the same partition. The default assertion policy is implementation-defined.
- 11/2 The following language-defined library package exists:
- ```
package Ada.Assertions is
  pragma Pure(Assertions);
  Assertion_Error : exception;
  procedure Assert(Check : in Boolean);
  procedure Assert(Check : in Boolean; Message : in String);
end Ada.Assertions;
```

A compilation unit containing a pragma Assert has a semantic dependence on the Assertions library unit.	16/2
The assertion policy that applies to a generic unit also applies to all its instances.	17/2

Dynamic Semantics

An assertion policy specifies how a **pragma** Assert is interpreted by the implementation. If the assertion policy is Ignore at the point of a **pragma** Assert, the pragma is ignored. If the assertion policy is Check at the point of a **pragma** Assert, the elaboration of the pragma consists of evaluating the boolean expression, and if the result is False, evaluating the Message argument, if any, and raising the exception `Assertions.Assertion_Error`, with a message if the Message argument is provided.

Calling the procedure `Assertions.Assert` without a Message parameter is equivalent to:

```
if Check = False then
  raise Ada.Assertions.Assertion_Error;
end if;
```

Calling the procedure `Assertions.Assert` with a Message parameter is equivalent to:

```
if Check = False then
  raise Ada.Assertions.Assertion_Error with Message;
end if;
```

The procedures `Assertions.Assert` have these effects independently of the assertion policy in effect.

Implementation Permissions

`Assertion_Error` may be declared by renaming an implementation-defined exception from another package.

Implementations may define their own assertion policies.

NOTES

2 Normally, the boolean expression in a **pragma** Assert should not call functions that have significant side-effects when the result of the expression is True, so that the particular assertion policy in effect will not affect normal operation of the program.

11.4.3 Example of Exception Handling

Examples

Exception handling may be used to separate the detection of an error from the response to that error:

```
package File_System is
  type File_Handle is limited private;
  File_Not_Found : exception;
  procedure Open(F : in out File_Handle; Name : String);
    -- raises File_Not_Found if named file does not exist
  End_Of_File : exception;
  procedure Read(F : in out File_Handle; Data : out Data_Type);
    -- raises End_Of_File if the file is not open
  ...
end File_System;
```

```

6/2      package body File_System is
7         procedure Open(F : in out File_Handle; Name : String) is
8             begin
9                 if File_Exists(Name) then
10                    ...
11                else
12                    raise File_Not_Found with "File not found: " & Name & ".";
13                end if;
14            end Open;
15
16            procedure Read(F : in out File_Handle; Data : out Data_Type) is
17                begin
18                    if F.Current_Position <= F.Last_Position then
19                        ...
20                    else
21                        raise End_Of_File;
22                    end if;
23                end Read;
24
25                ...
26
27            end File_System;
28
29            with Ada.Text_IO;
30            with Ada.Exceptions;
31            with File_System; use File_System;
32            use Ada;
33            procedure Main is
34                begin
35                    ... -- call operations in File_System
36                exception
37                    when End_Of_File =>
38                        Close(Some_File);
39                    when Not_Found_Error : File_Not_Found =>
40                        Text_IO.Put_Line(Exceptions.Exception_Message(Not_Found_Error));
41                    when The_Error : others =>
42                        Text_IO.Put_Line("Unknown error:");
43                        if Verbosity_Desired then
44                            Text_IO.Put_Line(Exceptions.Exception_Information(The_Error));
45                        else
46                            Text_IO.Put_Line(Exceptions.Exception_Name(The_Error));
47                            Text_IO.Put_Line(Exceptions.Exception_Message(The_Error));
48                        end if;
49                        raise;
50                end Main;

```

- 11 In the above example, the `File_System` package contains information about detecting certain exceptional situations, but it does not specify how to handle those situations. Procedure `Main` specifies how to handle them; other clients of `File_System` might have different handlers, even though the exceptional situations arise from the same basic causes.

11.5 Suppressing Checks

- 1/2 *Checking pragmas* give instructions to an implementation on handling language-defined checks. A `pragma Suppress` gives permission to an implementation to omit certain language-defined checks, while a `pragma Unsuppress` revokes the permission to omit checks..
- 2 A *language-defined check* (or simply, a “check”) is one of the situations defined by this International Standard that requires a check to be made at run time to determine whether some condition is true. A check *fails* when the condition being checked is false, causing an exception to be raised.

Syntax

The forms of checking pragmas are as follows:

pragma Suppress(identifier);

pragma Unsuppress(identifier);

A checking pragma is allowed only immediately within a **declarative_part**, immediately within a **package_specification**, or as a configuration pragma.

Legality Rules

The identifier shall be the name of a check.

This paragraph was deleted.

Static Semantics

A checking pragma applies to the named check in a specific region, and applies to all entities in that region. A checking pragma given in a **declarative_part** or immediately within a **package_specification** applies from the place of the **pragma** to the end of the innermost enclosing declarative region. The region for a checking pragma given as a configuration pragma is the declarative region for the entire compilation unit (or units) to which it applies.

If a checking pragma applies to a generic instantiation, then the checking pragma also applies to the instance. If a checking pragma applies to a call to a subprogram that has a **pragma** **Inline** applied to it, then the checking pragma also applies to the inlined subprogram body.

A **pragma** **Suppress** gives permission to an implementation to omit the named check (or every check in the case of **All_Checks**) for any entities to which it applies. If permission has been given to suppress a given check, the check is said to be *suppressed*.

A **pragma** **Unsuppress** revokes the permission to omit the named check (or every check in the case of **All_Checks**) given by any **pragma** **Suppress** that applies at the point of the **pragma** **Unsuppress**. The permission is revoked for the region to which the **pragma** **Unsuppress** applies. If there is no such permission at the point of a **pragma** **Unsuppress**, then the **pragma** has no effect. A later **pragma** **Suppress** can renew the permission.

The following are the language-defined checks:

- The following checks correspond to situations in which the exception **Constraint_Error** is raised upon failure.

Access_Check

When evaluating a dereference (explicit or implicit), check that the value of the name is not **null**. When converting to a subtype that excludes null, check that the converted value is not **null**.

Discriminant_Check

Check that the discriminants of a composite value have the values imposed by a discriminant constraint. Also, when accessing a record component, check that it exists for the current discriminant values.

Division_Check

Check that the second operand is not zero for the operations **/**, **rem** and **mod**.

Index_Check

Check that the bounds of an array value are equal to the corresponding bounds of an index constraint. Also, when accessing a component of an array object, check for each dimension that the given index value belongs to the range defined by the bounds of the

array object. Also, when accessing a slice of an array object, check that the given discrete range is compatible with the range defined by the bounds of the array object.

15 **Length_Check**

Check that two arrays have matching components, in the case of array subtype conversions, and logical operators for arrays of boolean components.

16 **Overflow_Check**

Check that a scalar value is within the base range of its type, in cases where the implementation chooses to raise an exception instead of returning the correct mathematical result.

17 **Range_Check**

Check that a scalar value satisfies a range constraint. Also, for the elaboration of a **subtype_indication**, check that the **constraint** (if present) is compatible with the subtype denoted by the **subtype_mark**. Also, for an **aggregate**, check that an index or discriminant value belongs to the corresponding subtype. Also, check that when the result of an operation yields an array, the value of each component belongs to the component subtype.

18 **Tag_Check**

Check that operand tags in a dispatching call are all equal. Check for the correct tag on tagged type conversions, for an **assignment_statement**, and when returning a tagged limited object from a function.

- 19 • The following checks correspond to situations in which the exception **Program_Error** is raised upon failure.

19.1/2 **Accessibility_Check**

Check the accessibility level of an entity or view.

19.2/2 **Allocation_Check**

For an **allocator**, check that the master of any tasks to be created by the **allocator** is not yet completed or some dependents have not yet terminated, and that the finalization of the collection has not started.

20 **Elaboration_Check**

When a subprogram or protected entry is called, a task activation is accomplished, or a generic instantiation is elaborated, check that the body of the corresponding unit has already been elaborated.

21/2 *This paragraph was deleted.*

- 22 • The following check corresponds to situations in which the exception **Storage_Error** is raised upon failure.

23 **Storage_Check**

Check that evaluation of an **allocator** does not require more space than is available for a storage pool. Check that the space available for a task or subprogram has not been exceeded.

- 24 • The following check corresponds to all situations in which any predefined exception is raised.

25 **All_Checks**

Represents the union of all checks; suppressing All_Checks suppresses all checks.

Erroneous Execution

- 26 If a given check has been suppressed, and the corresponding error situation occurs, the execution of the program is erroneous.

Implementation Permissions

An implementation is allowed to place restrictions on checking pragmas, subject only to the requirement that `pragma Unsuppress` shall allow any check names supported by `pragma Suppress`. An implementation is allowed to add additional check names, with implementation-defined semantics. When `Overflow_Check` has been suppressed, an implementation may also suppress an unspecified subset of the `Range_Checks`. 27/2

An implementation may support an additional parameter on `pragma Unsuppress` similar to the one allowed for `pragma Suppress` (see J.10). The meaning of such a parameter is implementation-defined. 27.1/2

Implementation Advice

The implementation should minimize the code executed for checks that have been suppressed. 28

NOTES

3 There is no guarantee that a suppressed check is actually removed; hence a `pragma Suppress` should be used only for efficiency reasons. 29

4 It is possible to give both a `pragma Suppress` and `Unsuppress` for the same check immediately within the same `declarative_part`. In that case, the last `pragma` given determines whether or not the check is suppressed. Similarly, it is possible to resuppress a check which has been unsuppressed by giving a `pragma Suppress` in an inner declarative region. 29.1/2

Examples

Examples of suppressing and unsuppressing checks: 30/2

```
pragma Suppress(Index_Check);
pragma Unsuppress(Overflow_Check);
```

31/2

11.6 Exceptions and Optimization

This clause gives permission to the implementation to perform certain “optimizations” that do not necessarily preserve the canonical semantics. 1

Dynamic Semantics

The rest of this International Standard (outside this clause) defines the *canonical semantics* of the language. The canonical semantics of a given (legal) program determines a set of possible external effects that can result from the execution of the program with given inputs. 2

As explained in 1.1.3, “Conformity of an Implementation with the Standard”, the external effect of a program is defined in terms of its interactions with its external environment. Hence, the implementation can perform any internal actions whatsoever, in any order or in parallel, so long as the external effect of the execution of the program is one that is allowed by the canonical semantics, or by the rules of this clause. 3

Implementation Permissions

The following additional permissions are granted to the implementation: 4

- An implementation need not always raise an exception when a language-defined check fails. Instead, the operation that failed the check can simply yield an *undefined result*. The exception need be raised by the implementation only if, in the absence of raising it, the value of this undefined result would have some effect on the external interactions of the program. In determining this, the implementation shall not presume that an undefined result has a value that belongs to its subtype, nor even to the base range of its type, if scalar. Having removed the raise of the exception, the canonical semantics will in general allow the implementation to omit the code for the check, and some or all of the operation itself. 5

- 6 • If an exception is raised due to the failure of a language-defined check, then upon reaching the corresponding `exception_handler` (or the termination of the task, if none), the external interactions that have occurred need reflect only that the exception was raised somewhere within the execution of the `sequence_of_statements` with the handler (or the `task_body`), possibly earlier (or later if the interactions are independent of the result of the checked operation) than that defined by the canonical semantics, but not within the execution of some abort-deferred operation or *independent* subprogram that does not dynamically enclose the execution of the construct whose check failed. An independent subprogram is one that is defined outside the library unit containing the construct whose check failed, and has no `Inline pragma` applied to it. Any assignment that occurred outside of such abort-deferred operations or independent subprograms can be disrupted by the raising of the exception, causing the object or its parts to become abnormal, and certain subsequent uses of the object to be erroneous, as explained in 13.9.1.

NOTES

- 7 5 The permissions granted by this clause can have an effect on the semantics of a program only if the program fails a language-defined check.

Section 12: Generic Units

A *generic unit* is a program unit that is either a generic subprogram or a generic package. A generic unit is a *template*, which can be parameterized, and from which corresponding (nongeneric) subprograms or packages can be obtained. The resulting program units are said to be *instances* of the original generic unit.

A generic unit is declared by a `generic_declaration`. This form of declaration has a `generic_formal_part` declaring any generic formal parameters. An instance of a generic unit is obtained as the result of a `generic_instantiation` with appropriate generic actual parameters for the generic formal parameters. An instance of a generic subprogram is a subprogram. An instance of a generic package is a package.

Generic units are templates. As templates they do not have the properties that are specific to their nongeneric counterparts. For example, a generic subprogram can be instantiated but it cannot be called. In contrast, an instance of a generic subprogram is a (nongeneric) subprogram; hence, this instance can be called but it cannot be used to produce further instances.

12.1 Generic Declarations

A `generic_declaration` declares a generic unit, which is either a generic subprogram or a generic package. A `generic_declaration` includes a `generic_formal_part` declaring any generic formal parameters. A generic formal parameter can be an object; alternatively (unlike a parameter of a subprogram), it can be a type, a subprogram, or a package.

Syntax

```

generic_declaration ::= generic_subprogram_declaration | generic_package_declaration
generic_subprogram_declaration ::= generic_formal_part subprogram_specification;
generic_package_declaration ::= generic_formal_part package_specification;
generic_formal_part ::= generic {generic_formal_parameter_declaration | use_clause}
generic_formal_parameter_declaration ::= formal_object_declaration
| formal_type_declaration
| formal_subprogram_declaration
| formal_package_declaration

```

The only form of `subtype_indication` allowed within a `generic_formal_part` is a `subtype_mark` (that is, the `subtype_indication` shall not include an explicit constraint). The defining name of a generic subprogram shall be an identifier (not an `operator_symbol`).

Static Semantics

A `generic_declaration` declares a generic unit — a generic package, generic procedure, or generic function, as appropriate.

An entity is a *generic formal* entity if it is declared by a `generic_formal_parameter_declaration`. “Generic formal,” or simply “formal,” is used as a prefix in referring to objects, subtypes (and types), functions, procedures and packages, that are generic formal entities, as well as to their respective declarations. Examples: “generic formal procedure” or a “formal integer type declaration.”

Dynamic Semantics

- ¹⁰ The elaboration of a generic declaration has no effect.

NOTES

- 1 Outside a generic unit a name that denotes the `generic_declaration` denotes the generic unit. In contrast, within the declarative region of the generic unit, a name that denotes the `generic_declaration` denotes the current instance.

- 2 Within a generic `subprogram_body`, the name of this program unit acts as the name of a subprogram. Hence this name can be overloaded, and it can appear in a recursive call of the current instance. For the same reason, this name cannot appear after the reserved word `new` in a (recursive) generic instantiation.

- 3 A `default_expression` or `default_name` appearing in a `generic_formal_part` is not evaluated during elaboration of the `generic_formal_part`; instead, it is evaluated when used. (The usual visibility rules apply to any `name` used in a default: the denoted declaration therefore has to be visible at the place of the expression.)

Examples

- #### 14 Examples of generic formal parts:

```

generic      -- parameterless
generic
  Size : Natural;   -- formal object
generic
  Length : Integer := 200;           -- formal object with a default expression
  Area    : Integer := Length*Length; -- formal object with a default expression
generic
  type Item  is private;           -- formal type
  type Index is (<>);           -- formal type
  type Row   is array(Index range <>) of Item; -- formal type
  with function "<"(X, Y : Item) return Boolean;  -- formal subprogram

```

- Examples of generic declarations declaring generic subprograms Exchange and Squaring:*

```
generic
    type Elemt is private;
procedure Exchange(U, V : in out Elemt);
generic
    type Itemt is private;
    with function "*" (U, V : Item) return Item is <>;
function Squaring(X : Item) return Item;
```

- 23 *Example of a generic declaration declaring a generic package:*

```

generic
  type Item    is private;
  type Vector is array (Positive range <>) of Item,
  with function Sum(X, Y : Item) return Item;
package On_Vectors is
  function Sum  (A, B : Vector) return Vector;
  function Sigma(A      : Vector) return Item;
  Length_Error : exception;
end On_Vectors;

```

12.2 Generic Bodies

The body of a generic unit (a *generic body*) is a template for the instance bodies. The syntax of a generic body is identical to that of a nongeneric body.

Dynamic Semantics

The elaboration of a generic body has no other effect than to establish that the generic unit can from then on be instantiated without failing the Elaboration_Check. If the generic body is a child of a generic package, then its elaboration establishes that each corresponding declaration nested in an instance of the parent (see 10.1.1) can from then on be instantiated without failing the Elaboration_Check.

NOTES

4 The syntax of generic subprograms implies that a generic subprogram body is always the completion of a declaration.

Example of a generic procedure body:

```
procedure Exchange(U, V : in out Elem) is -- see 12.1
    T : Elem; -- the generic formal type
begin
    T := U;
    U := V;
    V := T;
end Exchange;
```

Example of a generic function body:

```
function Squaring(X : Item) return Item is -- see 12.1
begin
    return X*X; -- the formal operator "*"
end Squaring;
```

Example of a generic package body:

```
package body On_Vectors is -- see 12.1
    function Sum(A, B : Vector) return Vector is
        Result : Vector(A'Range); -- the formal type Vector
        Bias   : constant Integer := B'First - A'First;
    begin
        if A'Length /= B'Length then
            raise Length_Error;
        end if;
        for N in A'Range loop
            Result(N) := Sum(A(N), B(N + Bias)); -- the formal function Sum
        end loop;
        return Result;
    end Sum;

    function Sigma(A : Vector) return Item is
        Total : Item := A(A'First); -- the formal type Item
    begin
        for N in A'First + 1 .. A'Last loop
            Total := Sum(Total, A(N)); -- the formal function Sum
        end loop;
        return Total;
    end Sigma;
end On_Vectors;
```

12.3 Generic Instantiation

- 1 An instance of a generic unit is declared by a generic_instantiation.

Syntax

```

2/2 generic_instantiation ::= 
    package defining_program_unit_name is
        new generic_package_name [generic_actual_part];
    | [overriding_indicator]
    | procedure defining_program_unit_name is
        new generic_procedure_name [generic_actual_part];
    | [overriding_indicator]
    | function defining_designator is
        new generic_function_name [generic_actual_part];
3 generic_actual_part ::= 
    (generic_association {, generic_association})
4 generic_association ::= 
    [generic_formal_parameter_selector_name =>] explicit_generic_actual_parameter
5 explicit_generic_actual_parameter ::= expression | variable_name
    | subprogram_name | entry_name | subtype_mark
    | package_instance_name
6 A generic_association is named or positional according to whether or not the generic_formal_
parameter_selector_name is specified. Any positional associations shall precede any named
associations.
7/2 The generic actual parameter is either the explicit_generic_actual_parameter given in a generic-
association for each formal, or the corresponding default_expression or default_name if no generic-
association is given for the formal. When the meaning is clear from context, the term “generic actual,” or
simply “actual,” is used as a synonym for “generic actual parameter” and also for the view denoted by one,
or the value of one.

```

Legality Rules

- 8 In a generic_instantiation for a particular kind of program unit (package, procedure, or function), the name shall denote a generic unit of the corresponding kind (generic package, generic procedure, or generic function, respectively).
- 9 The generic_formal_parameter_selector_name of a generic_association shall denote a generic_formal_parameter_declaration of the generic unit being instantiated. If two or more formal subprograms have the same defining name, then named associations are not allowed for the corresponding actuals.
- 10 A generic_instantiation shall contain at most one generic_association for each formal. Each formal without an association shall have a default_expression or subprogram_default.
- 11 In a generic unit Legality Rules are enforced at compile time of the generic_declaration and generic body, given the properties of the formals. In the visible part and formal part of an instance, Legality Rules are enforced at compile time of the generic_instantiation, given the properties of the actuals. In other parts of an instance, Legality Rules are not enforced; this rule does not apply when a given rule explicitly specifies otherwise.

Static Semantics

A `generic_instantiation` declares an instance; it is equivalent to the instance declaration (a `package_declaration` or `subprogram_declaration`) immediately followed by the instance body, both at the place of the instantiation.

The instance is a copy of the text of the template. Each use of a formal parameter becomes (in the copy) a use of the actual, as explained below. An instance of a generic package is a package, that of a generic procedure is a procedure, and that of a generic function is a function.

The interpretation of each construct within a generic declaration or body is determined using the overloading rules when that generic declaration or body is compiled. In an instance, the interpretation of each (copied) construct is the same, except in the case of a name that denotes the `generic_declaration` or some declaration within the generic unit; the corresponding name in the instance then denotes the corresponding copy of the denoted declaration. The overloading rules do not apply in the instance.

In an instance, a `generic_formal_parameter_declaration` declares a view whose properties are identical to those of the actual, except as specified in 12.4, “Formal Objects” and 12.6, “Formal Subprograms”. Similarly, for a declaration within a `generic_formal_parameter_declaration`, the corresponding declaration in an instance declares a view whose properties are identical to the corresponding declaration within the declaration of the actual.

Implicit declarations are also copied, and a name that denotes an implicit declaration in the generic denotes the corresponding copy in the instance. However, for a type declared within the visible part of the generic, a whole new set of primitive subprograms is implicitly declared for use outside the instance, and may differ from the copied set if the properties of the type in some way depend on the properties of some actual type specified in the instantiation. For example, if the type in the generic is derived from a formal private type, then in the instance the type will inherit subprograms from the corresponding actual type.

These new implicit declarations occur immediately after the type declaration in the instance, and override the copied ones. The copied ones can be called only from within the instance; the new ones can be called only from outside the instance, although for tagged types, the body of a new one can be executed by a call to an old one.

In the visible part of an instance, an explicit declaration overrides an implicit declaration if they are homographs, as described in 8.3. On the other hand, an explicit declaration in the private part of an instance overrides an implicit declaration in the instance, only if the corresponding explicit declaration in the generic overrides a corresponding implicit declaration in the generic. Corresponding rules apply to the other kinds of overriding described in 8.3.

Post-Compilation Rules

Recursive generic instantiation is not allowed in the following sense: if a given generic unit includes an instantiation of a second generic unit, then the instance generated by this instantiation shall not include an instance of the first generic unit (whether this instance is generated directly, or indirectly by intermediate instantiations).

Dynamic Semantics

For the elaboration of a `generic_instantiation`, each `generic_association` is first evaluated. If a default is used, an implicit `generic_association` is assumed for this rule. These evaluations are done in an arbitrary order, except that the evaluation for a default actual takes place after the evaluation for another actual if the default includes a name that denotes the other one. Finally, the instance declaration and body are elaborated.

- 21 For the evaluation of a generic_association the generic actual parameter is evaluated. Additional actions are performed in the case of a formal object of mode **in** (see 12.4).

NOTES

22 5 If a formal type is not tagged, then the type is treated as an untagged type within the generic body. Deriving from such a type in a generic body is permitted; the new type does not get a new tag value, even if the actual is tagged. Overriding operations for such a derived type cannot be dispatched to from outside the instance.

Examples

- 23 Examples of generic instantiations (see 12.1):

```
24    procedure Swap is new Exchange(Elem => Integer);
procedure Swap is new Exchange(Character);           -- Swap is overloaded
function Square is new Squaring(Integer);           -- "*" of Integer used by default
function Square is new Squaring(Item => Matrix, "*" => Matrix_Product);
function Square is new Squaring(Matrix, Matrix_Product); -- same as previous
25    package Int_Vectors is new On_Vectors(Integer, Table, "+");
```

- 26 Examples of uses of instantiated units:

```
27    Swap(A, B);
A := Square(A);
28    T : Table(1 .. 5) := (10, 20, 30, 40, 50);
N : Integer := Int_Vectors.Sigma(T); -- 150 (see 12.2, "Generic Bodies" for the body of Sigma)
29    use Int_Vectors;
M : Integer := Sigma(T); -- 150
```

12.4 Formal Objects

- 1 A generic formal object can be used to pass a value or variable to a generic unit.

Syntax

```
2/2    formal_object_declaration ::=
        defining_identifier_list : mode [null_exclusion] subtype_mark [:default_expression];
        defining_identifier_list : mode access_definition [:default_expression];
```

Name Resolution Rules

- 3 The expected type for the **default_expression**, if any, of a formal object is the type of the formal object.
- 4 For a generic formal object of mode **in**, the expected type for the actual is the type of the formal.
- 5/2 For a generic formal object of mode **in out**, the type of the actual shall resolve to the type determined by the **subtype_mark**, or for a **formal_object_declaration** with an **access_definition**, to a specific anonymous access type. If the anonymous access type is an access-to-object type, the type of the actual shall have the same designated type as that of the **access_definition**. If the anonymous access type is an access-to-subprogram type, the type of the actual shall have a designated profile which is type conformant with that of the **access_definition**.

Legality Rules

- 6 If a generic formal object has a **default_expression**, then the mode shall be **in** (either explicitly or by default); otherwise, its mode shall be either **in** or **in out**.
- 7 For a generic formal object of mode **in**, the actual shall be an **expression**. For a generic formal object of mode **in out**, the actual shall be a name that denotes a variable for which renaming is allowed (see 8.5.1).

In the case where the type of the formal is defined by an `access_definition`, the type of the actual and the type of the formal: 8/2

- shall both be access-to-object types with statically matching designated subtypes and with both or neither being access-to-constant types; or 8.1/2
- shall both be access-to-subprogram types with subtype conformant designated profiles. 8.2/2

For a `formal_object_declaration` with a `null_exclusion` or an `access_definition` that has a `null_exclusion`: 8.3/2

- if the actual matching the `formal_object_declaration` denotes the generic formal object of another generic unit G , and the instantiation containing the actual occurs within the body of G or within the body of a generic unit declared within the declarative region of G , then the declaration of the formal object of G shall have a `null_exclusion`; 8.4/2
- otherwise, the subtype of the actual matching the `formal_object_declaration` shall exclude null. In addition to the places where Legality Rules normally apply (see 12.3), this rule applies also in the private part of an instance of a generic unit. 8.5/2

Static Semantics

A `formal_object_declaration` declares a generic formal object. The default mode is `in`. For a formal object of mode `in`, the nominal subtype is the one denoted by the `subtype_mark` or `access_definition` in the declaration of the formal. For a formal object of mode `in out`, its type is determined by the `subtype_mark` or `access_definition` in the declaration; its nominal subtype is nonstatic, even if the `subtype_mark` denotes a static subtype; for a composite type, its nominal subtype is unconstrained if the first subtype of the type is unconstrained, even if the `subtype_mark` denotes a constrained subtype. 9/2

In an instance, a `formal_object_declaration` of mode `in` is a *full constant declaration* and declares a new stand-alone constant object whose initialization expression is the actual, whereas a `formal_object_declaration` of mode `in out` declares a view whose properties are identical to those of the actual. 10/2

Dynamic Semantics

For the evaluation of a `generic_association` for a formal object of mode `in`, a constant object is created, the value of the actual parameter is converted to the nominal subtype of the formal object, and assigned to the object, including any value adjustment — see 7.6. 11

NOTES

6 The constraints that apply to a generic formal object of mode `in out` are those of the corresponding generic actual parameter (not those implied by the `subtype_mark` that appears in the `formal_object_declaration`). Therefore, to avoid confusion, it is recommended that the name of a first subtype be used for the declaration of such a formal object. 12

12.5 Formal Types

- 1/2 A generic formal subtype can be used to pass to a generic unit a subtype whose type is in a certain category of types.

Syntax

```

2   formal_type_declaration ::= 
      type defining_identifier[discriminant_part] is formal_type_definition;
3/2  formal_type_definition ::= 
      formal_private_type_definition
      | formal_derived_type_definition
      | formal_discrete_type_definition
      | formal_signed_integer_type_definition
      | formal_modular_type_definition
      | formal_floating_point_definition
      | formal_ordinary_fixed_point_definition
      | formal_decimal_fixed_point_definition
      | formal_array_type_definition
      | formal_access_type_definition
      | formal_interface_type_definition

```

Legality Rules

- 4 For a generic formal subtype, the actual shall be a **subtype_mark**; it denotes the *(generic) actual subtype*.

Static Semantics

- 5 A **formal_type_declaration** declares a *(generic) formal type*, and its first subtype, the *(generic) formal subtype*.
- 6/2 The form of a **formal_type_definition** *determines a category (of types)* to which the formal type belongs. For a **formal_private_type_definition** the reserved words **tagged** and **limited** indicate the category of types (see 12.5.1). For a **formal_derived_type_definition** the category of types is the derivation class rooted at the ancestor type. For other formal types, the name of the syntactic category indicates the category of types; a **formal_discrete_type_definition** defines a discrete type, and so on.

Legality Rules

- 7/2 The actual type shall be in the category determined for the formal.

Static Semantics

- 8/2 The formal type also belongs to each category that contains the determined category. The primitive subprograms of the type are as for any type in the determined category. For a formal type other than a formal derived type, these are the predefined operators of the type. For an elementary formal type, the predefined operators are implicitly declared immediately after the declaration of the formal type. For a composite formal type, the predefined operators are implicitly declared either immediately after the declaration of the formal type, or later immediately within the declarative region in which the type is declared according to the rules of 7.3.1. In an instance, the copy of such an implicit declaration declares a view of the predefined operator of the actual type, even if this operator has been overridden for the actual type. The rules specific to formal derived types are given in 12.5.1.

NOTES

7 Generic formal types, like all types, are not named. Instead, a name can denote a generic formal subtype. Within a generic unit, a generic formal type is considered as being distinct from all other (formal or nonformal) types. 9

8 A discriminant_part is allowed only for certain kinds of types, and therefore only for certain kinds of generic formal types. See 3.7. 10

Examples

Examples of generic formal types:

```
type Item is private;
type Buffer(Length : Natural) is limited private;
type Enum is (<>);
type Int is range <>;
type Angle is delta <>;
type Mass is digits <>;
type Table is array (Enum) of Item;
```

Example of a generic formal part declaring a formal integer type:

```
generic
  type Rank is range <>;
  First : Rank := Rank'First;
  Second : Rank := First + 1; -- the operator "+" of the type Rank
```

12.5.1 Formal Private and Derived Types

In its most general form, the category determined for a formal private type is all types, but it can be restricted to only nonlimited types or to only tagged types. The category determined for a formal derived type is the derivation class rooted at the ancestor type. 1/2

Syntax

formal_private_type_definition ::= [[abstract] tagged] [limited] private 2

formal_derived_type_definition ::= 3/2
 [abstract] [limited | synchronized] new subtype_mark [[and interface_list]with private]

Legality Rules

If a generic formal type declaration has a known_discriminant_part, then it shall not include a default_expression for a discriminant. 4

The *ancestor subtype* of a formal derived type is the subtype denoted by the subtype_mark of the formal_derived_type_definition. For a formal derived type declaration, the reserved words with private shall appear if and only if the ancestor type is a tagged type; in this case the formal derived type is a private extension of the ancestor type and the ancestor shall not be a class-wide type. Similarly, an interface_list or the optional reserved words abstract or synchronized shall appear only if the ancestor type is a tagged type. The reserved word limited or synchronized shall appear only if the ancestor type and any progenitor types are limited types. The reserved word synchronized shall appear (rather than limited) if the ancestor type or any of the progenitor types are synchronized interfaces. 5/2

The actual type for a formal derived type shall be a descendant of the ancestor type and every progenitor of the formal type. If the reserved word synchronized appears in the declaration of the formal derived type, the actual type shall be a synchronized tagged type. 5.1/2

If the formal subtype is definite, then the actual subtype shall also be definite. 6

For a generic formal derived type with no discriminant_part: 7

- 8 • If the ancestor subtype is constrained, the actual subtype shall be constrained, and shall be statically compatible with the ancestor;
 - 9 • If the ancestor subtype is an unconstrained access or composite subtype, the actual subtype shall be unconstrained.
 - 10 • If the ancestor subtype is an unconstrained discriminated subtype, then the actual shall have the same number of discriminants, and each discriminant of the actual shall correspond to a discriminant of the ancestor, in the sense of 3.7.
 - 10.1/2 • If the ancestor subtype is an access subtype, the actual subtype shall exclude null if and only if the ancestor subtype excludes null.
- 11 The declaration of a formal derived type shall not have a `known_discriminant_part`. For a generic formal private type with a `known_discriminant_part`:
- 12 • The actual type shall be a type with the same number of discriminants.
 - 13 • The actual subtype shall be unconstrained.
 - 14 • The subtype of each discriminant of the actual type shall statically match the subtype of the corresponding discriminant of the formal type.
- 15 For a generic formal type with an `unknown_discriminant_part`, the actual may, but need not, have discriminants, and may be definite or indefinite.

Static Semantics

- 16/2 The category determined for a formal private type is as follows:
- | 17/2 <i>Type Definition</i> | Determined Category |
|-------------------------------|---|
| limited private | the category of all types |
| private | the category of all nonlimited types |
| tagged limited private | the category of all tagged types |
| tagged private | the category of all nonlimited tagged types |
- 18 The presence of the reserved word **abstract** determines whether the actual type may be abstract.
- 19 A formal private or derived type is a private or derived type, respectively. A formal derived tagged type is a private extension. A formal private or derived type is abstract if the reserved word **abstract** appears in its declaration.
- 20/2 If the ancestor type is a composite type that is not an array type, the formal type inherits components from the ancestor type (including discriminants if a new `discriminant_part` is not specified), as for a derived type defined by a `derived_type_definition` (see 3.4 and 7.3.1).
- 21/2 For a formal derived type, the predefined operators and inherited user-defined subprograms are determined by the ancestor type and any progenitor types, and are implicitly declared at the earliest place, if any, immediately within the declarative region in which the formal type is declared, where the corresponding primitive subprogram of the ancestor or progenitor is visible (see 7.3.1). In an instance, the copy of such an implicit declaration declares a view of the corresponding primitive subprogram of the ancestor or progenitor of the formal derived type, even if this primitive has been overridden for the actual type. When the ancestor or progenitor of the formal derived type is itself a formal type, the copy of the implicit declaration declares a view of the corresponding copied operation of the ancestor or progenitor. In the case of a formal private extension, however, the tag of the formal type is that of the actual type, so if the tag in a call is statically determined to be that of the formal type, the body executed will be that corresponding to the actual type.

For a prefix S that denotes a formal indefinite subtype, the following attribute is defined:

S'Definite S'Definite yields True if the actual subtype corresponding to S is definite; otherwise it yields False. The value of this attribute is of the predefined type Boolean.

Dynamic Semantics

In the case where a formal type is tagged with unknown discriminants, and the actual type is a class-wide type T'Class:

- For the purposes of defining the primitive operations of the formal type, each of the primitive operations of the actual type is considered to be a subprogram (with an intrinsic calling convention — see 6.3.1) whose body consists of a dispatching call upon the corresponding operation of T, with its formal parameters as the actual parameters. If it is a function, the result of the dispatching call is returned.
- If the corresponding operation of T has no controlling formal parameters, then the controlling tag value is determined by the context of the call, according to the rules for tag-indeterminate calls (see 3.9.2 and 5.2). In the case where the tag would be statically determined to be that of the formal type, the call raises Program_Error. If such a function is renamed, any call on the renaming raises Program_Error.

NOTES

9 In accordance with the general rule that the actual type shall belong to the category determined for the formal (see 12.5, “Formal Types”):

- If the formal type is nonlimited, then so shall be the actual;
- For a formal derived type, the actual shall be in the class rooted at the ancestor subtype.

10 The actual type can be abstract only if the formal type is abstract (see 3.9.3).

11 If the formal has a discriminant_part, the actual can be either definite or indefinite. Otherwise, the actual has to be definite.

12.5.2 Formal Scalar Types

A *formal scalar type* is one defined by any of the formal_type_definitions in this subclause. The category determined for a formal scalar type is the category of all discrete, signed integer, modular, floating point, ordinary fixed point, or decimal types.

Syntax

```
formal_discrete_type_definition ::= (<>)
formal_signed_integer_type_definition ::= range <>
formal_modular_type_definition ::= mod <>
formal_floating_point_definition ::= digits <>
formal_ordinary_fixed_point_definition ::= delta <>
formal_decimal_fixed_point_definition ::= delta <> digits <>
```

Legality Rules

The actual type for a formal scalar type shall not be a nonstandard numeric type.

NOTES

12 The actual type shall be in the class of types implied by the syntactic category of the formal type definition (see 12.5, “Formal Types”). For example, the actual for a formal_modular_type_definition shall be a modular type.

12.5.3 Formal Array Types

1/2 The category determined for a formal array type is the category of all array types.

Syntax

2 `formal_array_type_definition ::= array_type_definition`

Legality Rules

3 The only form of `discrete_subtype_definition` that is allowed within the declaration of a generic formal (constrained) array subtype is a `subtype_mark`.

4 For a formal array subtype, the actual subtype shall satisfy the following conditions:

- 5 • The formal array type and the actual array type shall have the same dimensionality; the formal subtype and the actual subtype shall be either both constrained or both unconstrained.
- 6 • For each index position, the index types shall be the same, and the index subtypes (if unconstrained), or the index ranges (if constrained), shall statically match (see 4.9.1).
- 7 • The component subtypes of the formal and actual array types shall statically match.
- 8 • If the formal type has aliased components, then so shall the actual.

Examples

9 Example of formal array types:

```
10      -- given the generic package
11
12      generic
13          type Item    is private;
14          type Index   is (<>);
15          type Vector  is array (Index range <>) of Item;
16          type Table   is array (Index) of Item;
17      package P is
18          .
19      end P;
20
21      -- and the types
22      type Mix     is array (Color range <>) of Boolean;
23      type Option  is array (Color) of Boolean;
24
25      -- then Mix can match Vector and Option can match Table
26      package R is new P(Item  => Boolean, Index => Color,
27                           Vector => Mix,           Table => Option);
28
29      -- Note that Mix cannot match Table and Option cannot match Vector
```

12.5.4 Formal Access Types

The category determined for a formal access type is the category of all access types.

1/2

Syntax

```
formal_access_type_definition ::= access_type_definition
```

2

Legality Rules

For a formal access-to-object type, the designated subtypes of the formal and actual types shall statically match.

3

If and only if the `general_access_modifier constant` applies to the formal, the actual shall be an access-to-constant type. If the `general_access_modifier all` applies to the formal, then the actual shall be a general access-to-variable type (see 3.10). If and only if the formal subtype excludes null, the actual subtype shall exclude null.

4/2

For a formal access-to-subprogram subtype, the designated profiles of the formal and the actual shall be mode-conformant, and the calling convention of the actual shall be *protected* if and only if that of the formal is *protected*.

5

Examples

Example of formal access types:

6

```
-- the formal types of the generic package
generic
  type Node is private;
  type Link is access Node;
  package P is
    ...
  end P;
-- can be matched by the actual types
type Car;
type Car_Name is access Car;
type Car is
  record
    Pred, Succ : Car_Name;
    Number     : License_Number;
    Owner      : Person;
  end record;
-- in the following generic instantiation
package R is new P(Node => Car, Link => Car_Name);
```

7

8

9

10

11

12

13

12.5.5 Formal Interface Types

The category determined for a formal interface type is the category of all interface types.

1/2

Syntax

```
formal_interface_type_definition ::= interface_type_definition
```

2/2

Legality Rules

The actual type shall be a descendant of every progenitor of the formal type.

3/2

- 4/2 The actual type shall be a limited, task, protected, or synchronized interface if and only if the formal type is also, respectively, a limited, task, protected, or synchronized interface.

Examples

```
5/2      type Root_Work_Item is tagged private;
6/2      generic
6/2          type Managed_Task is task interface;
6/2          type Work_Item(<>) is new Root_Work_Item with private;
6/2      package Server_Manager is
6/2          task type Server is new Managed_Task with
6/2              entry Start(Data : in out Work_Item);
6/2          end Server;
6/2      end Server_Manager;
```

- 7/2 This generic allows an application to establish a standard interface that all tasks need to implement so they can be managed appropriately by an application-specific scheduler.

12.6 Formal Subprograms

- 1 Formal subprograms can be used to pass callable entities to a generic unit.

Syntax

```
2/2      formal_subprogram_declaration ::= formal_concrete_subprogram_declaration
2/2          | formal_abstract_subprogram_declaration
2.1/2    formal_concrete_subprogram_declaration ::= 
2.1/2        with subprogram_specification [is subprogram_default];
2.2/2    formal_abstract_subprogram_declaration ::= 
2.2/2        with subprogram_specification is abstract [subprogram_default];
3/2      subprogram_default ::= default_name | <> | null
4       default_name ::= name
4.1/2    A subprogram_default of null shall not be specified for a formal function or for a
4.1/2      formal_abstract_subprogram_declaration.
```

Name Resolution Rules

- 5 The expected profile for the `default_name`, if any, is that of the formal subprogram.
- 6 For a generic formal subprogram, the expected profile for the actual is that of the formal subprogram.

Legality Rules

- 7 The profiles of the formal and any named default shall be mode-conformant.
- 8 The profiles of the formal and actual shall be mode-conformant.
- 8.1/2 For a parameter or result subtype of a `formal_subprogram_declaration` that has an explicit `null_exclusion`:
- 8.2/2 • if the actual matching the `formal_subprogram_declaration` denotes a generic formal object of another generic unit G , and the instantiation containing the actual that occurs within the body of a generic unit G or within the body of a generic unit declared within the declarative region of the generic unit G , then the corresponding parameter or result type of the formal subprogram of G shall have a `null_exclusion`;
- 8.3/2 • otherwise, the subtype of the corresponding parameter or result type of the actual matching the `formal_subprogram_declaration` shall exclude null. In addition to the places where Legality

Rules normally apply (see 12.3), this rule applies also in the private part of an instance of a generic unit.

If a formal parameter of a `formal_abstract_subprogram_declaration` is of a specific tagged type T or of an anonymous access type designating a specific tagged type T , T is called a *controlling type* of the `formal_abstract_subprogram_declaration`. Similarly, if the result of a `formal_abstract_subprogram_declaration` for a function is of a specific tagged type T or of an anonymous access type designating a specific tagged type T , T is called a controlling type of the `formal_abstract_subprogram_declaration`. A `formal_abstract_subprogram_declaration` shall have exactly one controlling type. 8.4/2

The actual subprogram for a `formal_abstract_subprogram_declaration` shall be a dispatching operation of the controlling type or of the actual type corresponding to the controlling type. 8.5/2

Static Semantics

A `formal_subprogram_declaration` declares a generic formal subprogram. The types of the formal parameters and result, if any, of the formal subprogram are those determined by the `subtype_marks` given in the `formal_subprogram_declaration`; however, independent of the particular subtypes that are denoted by the `subtype_marks`, the nominal subtypes of the formal parameters and result, if any, are defined to be nonstatic, and unconstrained if of an array type (no applicable index constraint is provided in a call on a formal subprogram). In an instance, a `formal_subprogram_declaration` declares a view of the actual. The profile of this view takes its subtypes and calling convention from the original profile of the actual entity, while taking the formal parameter names and `default_expressions` from the profile given in the `formal_subprogram_declaration`. The view is a function or procedure, never an entry. 9

If a generic unit has a `subprogram_default` specified by a box, and the corresponding actual parameter is omitted, then it is equivalent to an explicit actual parameter that is a usage name identical to the defining name of the formal. 10

If a generic unit has a `subprogram_default` specified by the reserved word `null`, and the corresponding actual parameter is omitted, then it is equivalent to an explicit actual parameter that is a null procedure having the profile given in the `formal_subprogram_declaration`. 10.1/2

The subprogram declared by a `formal_abstract_subprogram_declaration` with a controlling type T is a dispatching operation of type T . 10.2/2

NOTES

13 The matching rules for formal subprograms state requirements that are similar to those applying to `subprogram_renaming_declarations` (see 8.5.4). In particular, the name of a parameter of the formal subprogram need not be the same as that of the corresponding parameter of the actual subprogram; similarly, for these parameters, `default_expressions` need not correspond. 11

14 The constraints that apply to a parameter of a formal subprogram are those of the corresponding formal parameter of the matching actual subprogram (not those implied by the corresponding `subtype_mark` in the `_specification` of the formal subprogram). A similar remark applies to the result of a function. Therefore, to avoid confusion, it is recommended that the name of a first subtype be used in any declaration of a formal subprogram. 12

15 The subtype specified for a formal parameter of a generic formal subprogram can be any visible subtype, including a generic formal subtype of the same `generic_formal_part`. 13

16 A formal subprogram is matched by an attribute of a type if the attribute is a function with a matching specification. An enumeration literal of a given type matches a parameterless formal function whose result type is the given type. 14

17 A `default_name` denotes an entity that is visible or directly visible at the place of the `generic_declaration`; a box used as a default is equivalent to a name that denotes an entity that is directly visible at the place of the `_instantiation`. 15

18 The actual subprogram cannot be abstract unless the formal subprogram is a `formal_abstract_subprogram_declaration` (see 3.9.3). 16/2

- 16.1/2 19 The subprogram declared by a formal_abstract_subprogram_declaration is an abstract subprogram. All calls on a subprogram declared by a formal_abstract_subprogram_declaration must be dispatching calls. See 3.9.3.
- 16.2/2 20 A null procedure as a subprogram default has convention Intrinsic (see 6.3.1).

Examples

17 Examples of generic formal subprograms:

```
18/2  with function "+"(X, Y : Item) return Item is <>;
with function Image(X : Enum) return String is Enum'Image;
with procedure Update is Default_Update;
with procedure Pre_Action(X : in Item) is null; -- defaults to no action
with procedure Write(S : not null access Root_Stream_Type'Class;
                     Desc : Descriptor)
                     is abstract Descriptor'Write; -- see 13.13.2
-- Dispatching operation on Descriptor with default
-- given the generic procedure declaration
20  generic
      with procedure Action (X : in Item);
procedure Iterate(Seq : in Item_Sequence);
-- and the procedure
22  procedure Put_Item(X : in Item);
-- the following instantiation is possible
23  procedure Put_List is new Iterate(Action => Put_Item);
24
```

12.7 Formal Packages

- 1 Formal packages can be used to pass packages to a generic unit. The formal_package_declaration declares that the formal package is an instance of a given generic package. Upon instantiation, the actual package has to be an instance of that generic package.

Syntax

```
2  formal_package_declaration ::= 
      with package defining_identifier is new generic_package_name formal_package_actual_part;
3/2  formal_package_actual_part ::= 
      ([others =>] <>)
      | [generic_actual_part]
      | (formal_package_association {, formal_package_association} [, others => <>])
3.1/2 formal_package_association ::= 
      generic_association
      | generic_formal_parameter_selector_name => <>
3.2/2 Any positional formal_package_associations shall precede any named
      formal_package_associations.
```

Legality Rules

- 4 The generic_package_name shall denote a generic package (the template for the formal package); the formal package is an instance of the template.
- 4.1/2 A formal_package_actual_part shall contain at most one formal_package_association for each formal parameter. If the formal_package_actual_part does not include "others => <>", each formal parameter without an association shall have a default_expression or subprogram_default.

The actual shall be an instance of the template. If the formal_package_actual_part is ($\langle\rangle$) or (others \Rightarrow $\langle\rangle$), then the actual may be any instance of the template; otherwise, certain of the actual parameters of the actual instance shall match the corresponding actual parameters of the formal package, determined as follows:

- If the formal_package_actual_part includes generic_associations as well as associations with $\langle\rangle$, then only the actual parameters specified explicitly with generic_associations are required to match; 5.1/2
- Otherwise, all actual parameters shall match, whether any actual parameter is given explicitly or by default. 5.2/2

The rules for matching of actual parameters between the actual instance and the formal package are as follows: 5.3/2

- For a formal object of mode in, the actuals match if they are static expressions with the same value, or if they statically denote the same constant, or if they are both the literal null. 6/2
- For a formal subtype, the actuals match if they denote statically matching subtypes. 7
- For other kinds of formals, the actuals match if they statically denote the same entity. 8

For the purposes of matching, any actual parameter that is the name of a formal object of mode in is replaced by the formal object's actual expression (recursively). 8.1/1

Static Semantics

A formal_package_declaration declares a generic formal package. 9

The visible part of a formal package includes the first list of basic_declarative_items of the package_specification. In addition, for each actual parameter that is not required to match, a copy of the declaration of the corresponding formal parameter of the template is included in the visible part of the formal package. If the copied declaration is for a formal type, copies of the implicit declarations of the primitive subprograms of the formal type are also included in the visible part of the formal package. 10/2

For the purposes of matching, if the actual instance A is itself a formal package, then the actual parameters of A are those specified explicitly or implicitly in the formal_package_actual_part for A, plus, for those not specified, the copies of the formal parameters of the template included in the visible part of A. 11/2

Examples

Example of a generic package with formal package parameters: 12/2

```
with Ada.Containers.Ordered_Maps; -- see A.18.6
generic
  with package Mapping_1 is new Ada.Containers.Ordered_Maps(<>);
  with package Mapping_2 is new Ada.Containers.Ordered_Maps
    (Key_Type => Mapping_1.Element_Type,
     others => <>);

  package Ordered_Join is
    -- Provide a "join" between two mappings
    subtype Key_Type is Mapping_1.Key_Type;
    subtype Element_Type is Mapping_2.Element_Type;
    function Lookup(Key : Key_Type) return Element_Type;
    ...
  end Ordered_Join;
```

17/2 Example of an instantiation of a package with formal packages:

```

18/2   with Ada.Containers.Ordered_Maps;
19/2   package Symbol_Package is
20/2     type String_Id is ...
21/2     type Symbol_Info is ...
22/2   package String_Table is new Ada.Containers.Ordered_Maps
23/2     (Key_Type => String,
24/2       Element_Type => String_Id);
22/2   package Symbol_Table is new Ada.Containers.Ordered_Maps
23/2     (Key_Type => String_Id,
24/2       Element_Type => Symbol_Info);
23/2   package String_Info is new Ordered_Join(Mapping_1 => String_Table,
24/2                                         Mapping_2 => Symbol_Table);
24/2     Apple_Info : constant Symbol_Info := String_Info.Lookup("Apple");
25/2 end Symbol_Package;
```

12.8 Example of a Generic Package

1 The following example provides a possible formulation of stacks by means of a generic package. The size
of each stack and the type of the stack elements are provided as generic formal parameters.

Examples

2/1 This paragraph was deleted.

```

3   generic
4     Size : Positive;
5     type Item is private;
6     package Stack is
7       procedure Push(E : in Item);
8       procedure Pop (E : out Item);
9       Overflow, Underflow : exception;
10      end Stack;
11
12      package body Stack is
13        type Table is array (Positive range <>) of Item;
14        Space : Table(1 .. Size);
15        Index : Natural := 0;
16
17        procedure Push(E : in Item) is
18          begin
19            if Index >= Size then
20              raise Overflow;
21            end if;
22            Index := Index + 1;
23            Space(Index) := E;
24          end Push;
25
26        procedure Pop(E : out Item) is
27          begin
28            if Index = 0 then
29              raise Underflow;
30            end if;
31            E := Space(Index);
32            Index := Index - 1;
33          end Pop;
34
35      end Stack;
```

Instances of this generic package can be obtained as follows:

```
9
package Stack_Int  is new Stack(Size => 200, Item => Integer);
10 package Stack_Bool is new Stack(100, Boolean);
```

Thereafter, the procedures of the instantiated packages can be called as follows:

```
11
Stack_Int.Push(N);
12 Stack_Bool.Push(True);
```

Alternatively, a generic formulation of the type Stack can be given as follows (package body omitted):

```
13
generic
  type Item is private;
14 package On_Stacks is
  type Stack(Size : Positive) is limited private;
  procedure Push(S : in out Stack; E : in Item);
  procedure Pop (S : in out Stack; E : out Item);
  Overflow, Underflow : exception;
private
  type Table is array (Positive range <>) of Item;
  type Stack(Size : Positive) is
    record
      Space : Table(1 .. Size);
      Index : Natural := 0;
    end record;
end On_Stacks;
```

In order to use such a package, an instance has to be created and thereafter stacks of the corresponding type can be declared:

```
15
declare
  package Stack_Real is new On_Stacks(Real); use Stack_Real;
  S : Stack(100);
begin
  ...
  Push(S, 2.54);
  ...
end;
```

Section 13: Representation Issues

This section describes features for querying and controlling certain aspects of entities and for interfacing to hardware.

13.1 Operational and Representation Items

Representation and operational items can be used to specify aspects of entities. Two kinds of aspects of entities can be specified: aspects of representation and operational aspects. Representation items specify how the types and other entities of the language are to be mapped onto the underlying machine. Operational items specify other properties of entities.

There are six kinds of *representation items*: *attribute_definition_clauses* for representation attributes, *enumeration_representation_clauses*, *record_representation_clauses*, *at_clauses*, *component_clauses*, and *representation pragmas*. They can be provided to give more efficient representation or to interface with features that are outside the domain of the language (for example, peripheral hardware).

An *operational item* is an *attribute_definition_clause* for an operational attribute.

An operational item or a representation item applies to an entity identified by a *local_name*, which denotes an entity declared local to the current declarative region, or a library unit declared immediately preceding a representation pragma in a compilation.

Syntax

```
aspect_clause ::= attribute_definition_clause
  | enumeration_representation_clause
  | record_representation_clause
  | at_clause

local_name ::= direct_name
  | direct_name'attribute_designator
  | library_unit_name
```

A representation pragma is allowed only at places where an *aspect_clause* or *compilation_unit* is allowed.

Name Resolution Rules

In an operational item or representation item, if the *local_name* is a *direct_name*, then it shall resolve to denote a declaration (or, in the case of a *pragma*, one or more declarations) that occurs immediately within the same declarative region as the item. If the *local_name* has an *attribute_designator*, then it shall resolve to denote an implementation-defined component (see 13.5.1) or a class-wide type implicitly declared immediately within the same declarative region as the item. A *local_name* that is a *library_unit_name* (only permitted in a representation pragma) shall resolve to denote the *library_item* that immediately precedes (except for other pragmas) the representation pragma.

Legality Rules

The *local_name* of an *aspect_clause* or representation pragma shall statically denote an entity (or, in the case of a *pragma*, one or more entities) declared immediately preceding it in a *compilation*, or within the same *declarative_part*, *package_specification*, *task_definition*, *protected_definition*, or *record_definition* as the representation or operational item. If a *local_name* denotes a local callable entity, it may do so

through a local `subprogram_renaming_declaration` (as a way to resolve ambiguity in the presence of overloading); otherwise, the `local_name` shall not denote a `renaming_declaration`.

- 7/2 The *representation* of an object consists of a certain number of bits (the *size* of the object). For an object of an elementary type, these are the bits that are normally read or updated by the machine code when loading, storing, or operating-on the value of the object. For an object of a composite type, these are the bits reserved for this object, and include bits occupied by subcomponents of the object. If the size of an object is greater than that of its subtype, the additional bits are padding bits. For an elementary object, these padding bits are normally read and updated along with the others. For a composite object, padding bits might not be read or updated in any given composite operation, depending on the implementation.
- 8 A representation item *directly specifies* an *aspect of representation* of the entity denoted by the `local_name`, except in the case of a type-related representation item, whose `local_name` shall denote a first subtype, and which directly specifies an aspect of the subtype's type. A representation item that names a subtype is either *subtype-specific* (Size and Alignment clauses) or *type-related* (all others). Subtype-specific aspects may differ for different subtypes of the same type.
- 8.1/1 An operational item *directly specifies* an *operational aspect* of the type of the subtype denoted by the `local_name`. The `local_name` of an operational item shall denote a first subtype. An operational item that names a subtype is type-related.
- 9 A representation item that directly specifies an aspect of a subtype or type shall appear after the type is completely defined (see 3.11.1), and before the subtype or type is frozen (see 13.14). If a representation item is given that directly specifies an aspect of an entity, then it is illegal to give another representation item that directly specifies the same aspect of the entity.
- 9.1/1 An operational item that directly specifies an aspect of a type shall appear before the type is frozen (see 13.14). If an operational item is given that directly specifies an aspect of a type, then it is illegal to give another operational item that directly specifies the same aspect of the type.
- 10 For an untagged derived type, no type-related representation items are allowed if the parent type is a by-reference type, or has any user-defined primitive subprograms.
- 11/2 Operational and representation aspects of a generic formal parameter are the same as those of the actual. Operational and representation aspects are the same for all views of a type. A type-related representation item is not allowed for a descendant of a generic formal untagged type.
- 12 A representation item that specifies the Size for a given subtype, or the size or storage place for an object (including a component) of a given subtype, shall allow for enough storage space to accommodate any value of the subtype.
- 13/1 A representation or operational item that is not supported by the implementation is illegal, or raises an exception at run time.
- 13.1/2 A `type_declaration` is illegal if it has one or more progenitors, and a representation item applies to an ancestor, and this representation item conflicts with the representation of some other ancestor. The cases that cause conflicts are implementation defined.

Static Semantics

- 14 If two subtypes statically match, then their subtype-specific aspects (Size and Alignment) are the same.
- 15/1 A derived type inherits each type-related aspect of representation of its parent type that was directly specified before the declaration of the derived type, or (in the case where the parent is derived) that was inherited by the parent type from the grandparent type. A derived subtype inherits each subtype-specific

aspect of representation of its parent subtype that was directly specified before the declaration of the derived type, or (in the case where the parent is derived) that was inherited by the parent subtype from the grandparent subtype, but only if the parent subtype statically matches the first subtype of the parent type. An inherited aspect of representation is overridden by a subsequent representation item that specifies the same aspect of the type or subtype.

In contrast, whether operational aspects are inherited by an untagged derived type depends on each specific aspect. Operational aspects are never inherited for a tagged type. When operational aspects are inherited by an untagged derived type, aspects that were directly specified by operational items that are visible at the point of the derived type declaration, or (in the case where the parent is derived) that were inherited by the parent type from the grandparent type are inherited. An inherited operational aspect is overridden by a subsequent operational item that specifies the same aspect of the type. 15.1/2

When an aspect that is a subprogram is inherited, the derived type inherits the aspect in the same way that a derived type inherits a user-defined primitive subprogram from its parent (see 3.4). 15.2/2

Each aspect of representation of an entity is as follows: 16

- If the aspect is *specified* for the entity, meaning that it is either directly specified or inherited, then that aspect of the entity is as specified, except in the case of Storage_Size, which specifies a minimum. 17
- If an aspect of representation of an entity is not specified, it is chosen by default in an unspecified manner. 18

If an operational aspect is *specified* for an entity (meaning that it is either directly specified or inherited), then that aspect of the entity is as specified. Otherwise, the aspect of the entity has the default value for that aspect. 18.1/1

A representation item that specifies an aspect of representation that would have been chosen in the absence of the representation item is said to be *confirming*. 18.2/2

Dynamic Semantics

For the elaboration of an `aspect_clause`, any evaluable constructs within it are evaluated. 19/1

Implementation Permissions

An implementation may interpret aspects of representation in an implementation-defined manner. An implementation may place implementation-defined restrictions on representation items. A *recommended level of support* is specified for representation items and related features in each subclause. These recommendations are changed to requirements for implementations that support the Systems Programming Annex (see C.2, “Required Representation Support”). 20

Implementation Advice

The recommended level of support for all representation items is qualified as follows: 21

- A confirming representation item should be supported. 21.1/2
- An implementation need not support representation items containing nonstatic expressions, except that an implementation should support a representation item for a given entity if each nonstatic expression in the representation item is a name that statically denotes a constant declared before the entity. 22
- An implementation need not support a specification for the Size for a given composite subtype, nor the size or storage place for an object (including a component) of a given composite subtype, 23

- unless the constraints on the subtype and its composite subcomponents (if any) are all static constraints.
- 24/2 • An implementation need not support a nonconfirming representation item if it could cause an aliased object or an object of a by-reference type to be allocated at a nonaddressable location or, when the alignment attribute of the subtype of such an object is nonzero, at an address that is not an integral multiple of that alignment.
 - 25/2 • An implementation need not support a nonconfirming representation item if it could cause an aliased object of an elementary type to have a size other than that which would have been chosen by default.
 - 26/2 • An implementation need not support a nonconfirming representation item if it could cause an aliased object of a composite type, or an object whose type is by-reference, to have a size smaller than that which would have been chosen by default.
 - 27/2 • An implementation need not support a nonconfirming subtype-specific representation item specifying an aspect of representation of an indefinite or abstract subtype.
- 28/2 For purposes of these rules, the determination of whether a representation item applied to a type *could cause* an object to have some property is based solely on the properties of the type itself, not on any available information about how the type is used. In particular, it presumes that minimally aligned objects of this type might be declared at some point.

13.2 Pragma Pack

- 1 A **pragma** Pack specifies that storage minimization should be the main criterion when selecting the representation of a composite type.

Syntax

- 2 The form of a **pragma** Pack is as follows:
- 3 **pragma** Pack(*first_subtype_local_name*);

Legality Rules

- 4 The *first_subtype_local_name* of a **pragma** Pack shall denote a composite subtype.

Static Semantics

- 5 A **pragma** Pack specifies the *packing* aspect of representation; the type (or the extension part) is said to be *packed*. For a type extension, the parent part is packed as for the parent type, and a **pragma** Pack causes packing only of the extension part.

Implementation Advice

- 6 If a type is packed, then the implementation should try to minimize storage allocated to objects of the type, possibly at the expense of speed of accessing components, subject to reasonable complexity in addressing calculations.
- 6.1/2 If a packed type has a component that is not of a by-reference type and has no aliased part, then such a component need not be aligned according to the Alignment of its subtype; in particular it need not be allocated on a storage element boundary.
- 7 The recommended level of support for **pragma** Pack is:
- 8 • For a packed record type, the components should be packed as tightly as possible subject to the Sizes of the component subtypes, and subject to any *record_representation_clause* that applies

to the type; the implementation may, but need not, reorder components or cross aligned word boundaries to improve the packing. A component whose Size is greater than the word size may be allocated an integral number of words.

- For a packed array type, if the component subtype's Size is less than or equal to the word size, and Component_Size is not specified for the type, Component_Size should be less than or equal to the Size of the component subtype, rounded up to the nearest factor of the word size.

9

13.3 Operational and Representation Attributes

The values of certain implementation-dependent characteristics can be obtained by interrogating appropriate operational or representation attributes. Some of these attributes are specifiable via an attribute_definition_clause.

1/1

Syntax

```
attribute_definition_clause ::=  
    for local_name'attribute_designator use expression;  
    | for local_name'attribute_designator use name;
```

2

Name Resolution Rules

For an attribute_definition_clause that specifies an attribute that denotes a value, the form with an expression shall be used. Otherwise, the form with a name shall be used.

3

For an attribute_definition_clause that specifies an attribute that denotes a value or an object, the expected type for the expression or name is that of the attribute. For an attribute_definition_clause that specifies an attribute that denotes a subprogram, the expected profile for the name is the profile required for the attribute. For an attribute_definition_clause that specifies an attribute that denotes some other kind of entity, the name shall resolve to denote an entity of the appropriate kind.

4

Legality Rules

An attribute_designator is allowed in an attribute_definition_clause only if this International Standard explicitly allows it, or for an implementation-defined attribute if the implementation allows it. Each specifiable attribute constitutes an operational aspect or aspect of representation.

5/1

For an attribute_definition_clause that specifies an attribute that denotes a subprogram, the profile shall be mode conformant with the one required for the attribute, and the convention shall be Ada. Additional requirements are defined for particular attributes.

6

Static Semantics

A *Size clause* is an attribute_definition_clause whose attribute_designator is Size. Similar definitions apply to the other specifiable attributes.

7/2

A *storage element* is an addressable element of storage in the machine. A *word* is the largest amount of storage that can be conveniently and efficiently manipulated by the hardware, given the implementation's run-time model. A word consists of an integral number of storage elements.

8

A *machine scalar* is an amount of storage that can be conveniently and efficiently loaded, stored, or operated upon by the hardware. Machine scalars consist of an integral number of storage elements. The set of machine scalars is implementation defined, but must include at least the storage element and the word. Machine scalars are used to interpret component_clauses when the nondefault bit ordering applies.

8.1/2

- 9/1 The following representation attributes are defined: Address, Alignment, Size, Storage_Size, and Component_Size.
- 10/1 For a prefix X that denotes an object, program unit, or label:
- 11 X'Address Denotes the address of the first of the storage elements allocated to X. For a program unit or label, this value refers to the machine code associated with the corresponding body or statement. The value of this attribute is of type System.Address.
- 12 Address may be specified for stand-alone objects and for program units via an attribute_definition_clause.

Erroneous Execution

- 13 If an Address is specified, it is the programmer's responsibility to ensure that the address is valid; otherwise, program execution is erroneous.

Implementation Advice

- 14 For an array X, X'Address should point at the first component of the array, and not at the array bounds.
- 15 The recommended level of support for the Address attribute is:
- 16 • X'Address should produce a useful result if X is an object that is aliased or of a by-reference type, or is an entity whose Address has been specified.
 - 17 • An implementation should support Address clauses for imported subprograms.
- 18/2 • *This paragraph was deleted.*
- 19 • If the Address of an object is specified, or it is imported or exported, then the implementation should not perform optimizations based on assumptions of no aliases.

NOTES

- 20 1 The specification of a link name in a pragma Export (see B.1) for a subprogram or object is an alternative to explicit specification of its link-time address, allowing a link-time directive to place the subprogram or object within memory.
- 21 2 The rules for the Size attribute imply, for an aliased object X, that if X'Size = Storage_Unit, then X'Address points at a storage element containing all of the bits of X, and only the bits of X.

Static Semantics

- 22/2 For a prefix X that denotes an object:
- 23/2 X'Alignment The value of this attribute is of type *universal_integer*, and nonnegative; zero means that the object is not necessarily aligned on a storage element boundary. If X'Alignment is not zero, then X is aligned on a storage unit boundary and X'Address is an integral multiple of X'Alignment (that is, the Address modulo the Alignment is zero).
- 24/2 *This paragraph was deleted.*
- 25/2 Alignment may be specified for stand-alone objects via an attribute_definition_clause; the expression of such a clause shall be static, and its value nonnegative.
- 26/2 *This paragraph was deleted.*
- 26.1/2 For every subtype S:
- 26.2/2 S'Alignment The value of this attribute is of type *universal_integer*, and nonnegative.
- 26.3/2 For an object X of subtype S, if S'Alignment is not zero, then X'Alignment is a nonzero integral multiple of S'Alignment unless specified otherwise by a representation item.
- 26.4/2 Alignment may be specified for first subtypes via an attribute_definition_clause; the expression of such a clause shall be static, and its value nonnegative.

Erroneous Execution

Program execution is erroneous if an Address clause is given that conflicts with the Alignment.	27
For an object that is not allocated under control of the implementation, execution is erroneous if the object is not aligned according to its Alignment.	28/2

Implementation Advice

The recommended level of support for the Alignment attribute for subtypes is:	29
<ul style="list-style-type: none"> An implementation should support an Alignment clause for a discrete type, fixed point type, record type, or array type, specifying an Alignment value that is zero or a power of two, subject to the following: 	30/2
<ul style="list-style-type: none"> An implementation need not support an Alignment clause for a signed integer type specifying an Alignment greater than the largest Alignment value that is ever chosen by default by the implementation for any signed integer type. A corresponding limitation may be imposed for modular integer types, fixed point types, enumeration types, record types, and array types. 	31/2
<ul style="list-style-type: none"> An implementation need not support a nonconfirming Alignment clause which could enable the creation of an object of an elementary type which cannot be easily loaded and stored by available machine instructions. 	32/2
<ul style="list-style-type: none"> An implementation need not support an Alignment specified for a derived tagged type which is not a multiple of the Alignment of the parent type. An implementation need not support a nonconfirming Alignment specified for a derived untagged by-reference type. 	32.1/2

The recommended level of support for the Alignment attribute for objects is:	33
<ul style="list-style-type: none"> <i>This paragraph was deleted.</i> 	34/2
<ul style="list-style-type: none"> For stand-alone library-level objects of statically constrained subtypes, the implementation should support all Alignments supported by the target linker. For example, page alignment is likely to be supported for such objects, but not for subtypes. 	35
<ul style="list-style-type: none"> For other objects, an implementation should at least support the alignments supported for their subtype, subject to the following: 	35.1/2
<ul style="list-style-type: none"> An implementation need not support Alignments specified for objects of a by-reference type or for objects of types containing aliased subcomponents if the specified Alignment is not a multiple of the Alignment of the subtype of the object. 	35.2/2

NOTES

3 Alignment is a subtype-specific attribute.	36
<i>This paragraph was deleted.</i>	37/2
4 A component_clause, Component_Size clause, or a pragma Pack can override a specified Alignment.	38

Static Semantics

For a prefix X that denotes an object:	39/1
X'Size Denotes the size in bits of the representation of the object. The value of this attribute is of the type <i>universal_integer</i> .	40
Size may be specified for stand-alone objects via an attribute_definition_clause; the expression of such a clause shall be static and its value nonnegative.	41

Implementation Advice

The size of an array object should not include its bounds.	41.1/2
--	--------

42/2 The recommended level of support for the Size attribute of objects is the same as for subtypes (see below), except that only a confirming Size clause need be supported for an aliased elementary object.

43/2 • This paragraph was deleted.

Static Semantics

44 For every subtype S:

45 S'Size If S is definite, denotes the size (in bits) that the implementation would choose for the following objects of subtype S:

- A record component of subtype S when the record type is packed.

46 47 • The formal parameter of an instance of `Unchecked_Conversion` that converts from subtype S to some other subtype.

48 If S is indefinite, the meaning is implementation defined. The value of this attribute is of the type `universal_integer`. The Size of an object is at least as large as that of its subtype, unless the object's Size is determined by a Size clause, a component_clause, or a Component_Size clause. Size may be specified for first subtypes via an attribute_definition_clause; the expression of such a clause shall be static and its value nonnegative.

Implementation Requirements

49 In an implementation, Boolean'Size shall be 1.

Implementation Advice

50/2 If the Size of a subtype allows for efficient independent addressability (see 9.10) on the target architecture, then the Size of the following objects of the subtype should equal the Size of the subtype:

- Aliased objects (including components).
- Unaliased components, unless the Size of the component is determined by a component_clause or Component_Size clause.

53 A Size clause on a composite subtype should not affect the internal layout of components.

54 The recommended level of support for the Size attribute of subtypes is:

- The Size (if not specified) of a static discrete or fixed point subtype should be the number of bits needed to represent each value belonging to the subtype using an unbiased representation, leaving space for a sign bit only if the subtype contains negative values. If such a subtype is a first subtype, then an implementation should support a specified Size for it that reflects this representation.
- For a subtype implemented with levels of indirection, the Size should include the size of the pointers, but not the size of what they point at.

56.1/2 • An implementation should support a Size clause for a discrete type, fixed point type, record type, or array type, subject to the following:

56.2/2 • An implementation need not support a Size clause for a signed integer type specifying a Size greater than that of the largest signed integer type supported by the implementation in the absence of a size clause (that is, when the size is chosen by default). A corresponding limitation may be imposed for modular integer types, fixed point types, enumeration types, record types, and array types.

56.3/2 • A nonconfirming size clause for the first subtype of a derived untagged by-reference type need not be supported.

NOTES

5 Size is a subtype-specific attribute.

57

6 A component_clause or Component_Size clause can override a specified Size. A pragma Pack cannot.

58

Static Semantics

For a prefix T that denotes a task object (after any implicit dereference):

59/1

T'Storage_Size

Denotes the number of storage elements reserved for the task. The value of this attribute is of the type *universal_integer*. The Storage_Size includes the size of the task's stack, if any. The language does not specify whether or not it includes other storage associated with the task (such as the “task control block” used by some implementations.) If a pragma Storage_Size is given, the value of the Storage_Size attribute is at least the value specified in the pragma.

60

A pragma Storage_Size specifies the amount of storage to be reserved for the execution of a task.

61

Syntax

The form of a pragma Storage_Size is as follows:

62

pragma Storage_Size(*expression*);

63

A pragma Storage_Size is allowed only immediately within a task_definition.

64

Name Resolution Rules

The expression of a pragma Storage_Size is expected to be of any integer type.

65

Dynamic Semantics

A pragma Storage_Size is elaborated when an object of the type defined by the immediately enclosing task_definition is created. For the elaboration of a pragma Storage_Size, the expression is evaluated; the Storage_Size attribute of the newly created task object is at least the value of the expression.

66

At the point of task object creation, or upon task activation, Storage_Error is raised if there is insufficient free storage to accommodate the requested Storage_Size.

67

Static Semantics

For a prefix X that denotes an array subtype or array object (after any implicit dereference):

68/1

X'Component_Size

Denotes the size in bits of components of the type of X. The value of this attribute is of type *universal_integer*.

69

Component_Size may be specified for array types via an attribute_definition_clause; the expression of such a clause shall be static, and its value nonnegative.

70

Implementation Advice

The recommended level of support for the Component_Size attribute is:

71

- An implementation need not support specified Component_Sizes that are less than the Size of the component subtype.
- An implementation should support specified Component_Sizes that are factors and multiples of the word size. For such Component_Sizes, the array should contain no gaps between components. For other Component_Sizes (if supported), the array should contain no gaps between components when packing is also specified; the implementation should forbid this combination in cases where it cannot support a no-gaps representation.

72

73

Static Semantics

- 73.1/1 The following operational attribute is defined: External_Tag.
- 74/1 For every subtype S of a tagged type T (specific or class-wide):
- 75/1 S'External_Tag

S'External_Tag denotes an external string representation for S'Tag; it is of the predefined type String. External_Tag may be specified for a specific tagged type via an attribute_definition_clause; the expression of such a clause shall be static. The default external tag representation is implementation defined. See 3.9.2 and 13.13.2. The value of External_Tag is never inherited; the default value is always used unless a new value is directly specified for a type.

Implementation Requirements

- 76 In an implementation, the default external tag for each specific tagged type declared in a partition shall be distinct, so long as the type is declared outside an instance of a generic body. If the compilation unit in which a given tagged type is declared, and all compilation units on which it semantically depends, are the same in two different partitions, then the external tag for the type shall be the same in the two partitions. What it means for a compilation unit to be the same in two different partitions is implementation defined. At a minimum, if the compilation unit is not recompiled between building the two different partitions that include it, the compilation unit is considered the same in the two partitions.

NOTES

- 77/2 7 The following language-defined attributes are specifiable, at least for some of the kinds of entities to which they apply: Address, Alignment, Bit_Order, Component_Size, External_Tag, Input, Machine_Radix, Output, Read, Size, Small, Storage_Pool, Storage_Size, Stream_Size, and Write.
- 78 8 It follows from the general rules in 13.1 that if one writes “**for X'Size use Y;**” then the X'Size attribute_reference will return Y (assuming the implementation allows the Size clause). The same is true for all of the specifiable attributes except Storage_Size.

Examples

- 79 Examples of attribute definition clauses:

```

80  Byte : constant := 8;
    Page : constant := 2**12;

81  type Medium is range 0 .. 65_000;
    for Medium'Size use 2*Byte;
    for Medium'Alignment use 2;
    Device_Register : Medium;
    for Device_Register'Size use Medium'Size;
    for Device_Register'Address use
        System.Storage_Elements.To_Address(16#FFFF_0020#);

82  type Short is delta 0.01 range -100.0 .. 100.0;
    for Short'Size use 15;

83  for Car_Name'Storage_Size use -- specify access type's storage pool size
      2000*((Car'Size/System.Storage_Unit) +1); -- approximately 2000 cars

84/2 function My_Input(Stream : not null access
    Ada.Streams.Root_Stream_Type'Class)
    return T;
    for T'Input use My_Input; -- see 13.13.2

```

NOTES

- 85 9 Notes on the examples: In the Size clause for Short, fifteen bits is the minimum necessary, since the type definition requires Short'Small <= 2**(-7).

13.4 Enumeration Representation Clauses

An enumeration_representation_clause specifies the internal codes for enumeration literals.

Syntax

```
enumeration_representation_clause ::=  
    for first_subtype_local_name use enumeration_aggregate;  
  
enumeration_aggregate ::= array_aggregate
```

Name Resolution Rules

The enumeration_aggregate shall be written as a one-dimensional array_aggregate, for which the index subtype is the unconstrained subtype of the enumeration type, and each component expression is expected to be of any integer type.

Legality Rules

The *first_subtype_local_name* of an enumeration_representation_clause shall denote an enumeration subtype.

Each component of the array_aggregate shall be given by an expression rather than a $\langle\rangle$. The expressions given in the array_aggregate shall be static, and shall specify distinct integer codes for each value of the enumeration type; the associated integer codes shall satisfy the predefined ordering relation of the type.

Static Semantics

An enumeration_representation_clause specifies the *coding* aspect of representation. The coding consists of the *internal code* for each enumeration literal, that is, the integral value used internally to represent each literal.

Implementation Requirements

For nonboolean enumeration types, if the coding is not specified for the type, then for each value of the type, the internal code shall be equal to its position number.

Implementation Advice

The recommended level of support for enumeration_representation_clauses is:

- An implementation should support at least the internal codes in the range System.Min_Int..System.Max_Int. An implementation need not support enumeration_representation_clauses for boolean types.

NOTES

10 Unchecked_Conversion may be used to query the internal codes used for an enumeration type. The attributes of the type, such as Succ, Pred, and Pos, are unaffected by the enumeration_representation_clause. For example, Pos always returns the position number, *not* the internal integer code that might have been specified in an enumeration_representation_clause}.

Examples

Example of an enumeration representation clause:

```
type Mix_Code is (ADD, SUB, MUL, LDA, STA, STZ);  
for Mix_Code use  
    (ADD => 1, SUB => 2, MUL => 3, LDA => 8, STA => 24, STZ => 33);
```

13.5 Record Layout

- 1 The (*record*) layout aspect of representation consists of the *storage places* for some or all components, that is, storage place attributes of the components. The layout can be specified with a `record_representation_-clause`.

13.5.1 Record Representation Clauses

- 1 A `record_representation_clause` specifies the storage representation of records and record extensions, that is, the order, position, and size of components (including discriminants, if any).

Syntax

```

2   record_representation_clause ::= 
3     for first_subtype_local_name use
4       record [mod_clause]
5         {component_clause}
6       end record;
7
8   component_clause ::= 
9     component_local_name at position range first_bit .. last_bit;
10
11  position ::= static_expression
12
13  first_bit ::= static_simple_expression
14
15  last_bit ::= static_simple_expression

```

Name Resolution Rules

- 7 Each `position`, `first_bit`, and `last_bit` is expected to be of any integer type.

Legality Rules

- 8/2 The `first_subtype_local_name` of a `record_representation_clause` shall denote a specific record or record extension subtype.
- 9 If the `component_local_name` is a `direct_name`, the `local_name` shall denote a component of the type. For a record extension, the component shall not be inherited, and shall not be a discriminant that corresponds to a discriminant of the parent type. If the `component_local_name` has an `attribute_designator`, the `direct_name` of the `local_name` shall denote either the declaration of the type or a component of the type, and the `attribute_designator` shall denote an implementation-defined implicit component of the type.
- 10 The `position`, `first_bit`, and `last_bit` shall be static expressions. The value of `position` and `first_bit` shall be nonnegative. The value of `last_bit` shall be no less than `first_bit - 1`.
- 10.1/2 If the nondefault bit ordering applies to the type, then either:
- 10.2/2 • the value of `last_bit` shall be less than the size of the largest machine scalar; or
 - 10.3/2 • the value of `first_bit` shall be zero and the value of `last_bit + 1` shall be a multiple of `System.Storage_Unit`.
- 11 At most one `component_clause` is allowed for each component of the type, including for each discriminant (`component_clauses` may be given for some, all, or none of the components). Storage places within a `component_list` shall not overlap, unless they are for components in distinct variants of the same `variant_part`.

A name that denotes a component of a type is not allowed within a `record_representation_clause` for the type, except as the `component_local_name` of a `component_clause`. 12

Static Semantics

A `record_representation_clause` (without the `mod_clause`) specifies the layout. 13/2

If the default bit ordering applies to the type, the `position`, `first_bit`, and `last_bit` of each `component_clause` directly specify the position and size of the corresponding component. 13.1/2

If the nondefault bit ordering applies to the type then the layout is determined as follows: 13.2/2

- the `component_clauses` for which the value of `last_bit` is greater than or equal to the size of the largest machine scalar directly specify the position and size of the corresponding component; 13.3/2
- for other `component_clauses`, all of the components having the same value of `position` are considered to be part of a single machine scalar, located at that `position`; this machine scalar has a size which is the smallest machine scalar size larger than the largest `last_bit` for all `component_clauses` at that `position`; the `first_bit` and `last_bit` of each `component_clause` are then interpreted as bit offsets in this machine scalar. 13.4/2

A `record_representation_clause` for a record extension does not override the layout of the parent part; if the layout was specified for the parent type, it is inherited by the record extension. 14

Implementation Permissions

An implementation may generate implementation-defined components (for example, one containing the offset of another component). An implementation may generate names that denote such implementation-defined components; such names shall be implementation-defined `attribute_references`. An implementation may allow such implementation-defined names to be used in `record_representation_clauses`. An implementation can restrict such `component_clauses` in any manner it sees fit. 15

If a `record_representation_clause` is given for an untagged derived type, the storage place attributes for all of the components of the derived type may differ from those of the corresponding components of the parent type, even for components whose storage place is not specified explicitly in the `record_representation_clause`. 16

Implementation Advice

The recommended level of support for `record_representation_clauses` is: 17

- An implementation should support machine scalars that correspond to all of the integer, floating point, and address formats supported by the machine. 17.1/2
- An implementation should support storage places that can be extracted with a load, mask, shift sequence of machine code, and set with a load, shift, mask, store sequence, given the available machine instructions and run-time model. 18
- A storage place should be supported if its size is equal to the `Size` of the component subtype, and it starts and ends on a boundary that obeys the `Alignment` of the component subtype. 19
- For a component with a subtype whose `Size` is less than the word size, any storage place that does not cross an aligned word boundary should be supported. 20/2
- An implementation may reserve a storage place for the tag field of a tagged type, and disallow other components from overlapping that place. 21
- An implementation need not support a `component_clause` for a component of an extension part if the storage place is not after the storage places of all components of the parent type, whether or not those storage places had been specified. 22

NOTES

23 11 If no `component_clause` is given for a component, then the choice of the storage place for the component is left to the implementation. If `component_clauses` are given for all components, the `record_representation_clause` completely specifies the representation of the type and will be obeyed exactly by the implementation.

Examples

24 *Example of specifying the layout of a record type:*

```
25     Word : constant := 4; -- storage element is byte, 4 bytes per word
26     type State      is (A,M,W,P);
27     type Mode       is (Fix, Dec, Exp, Signif);
28     type Byte_Mask   is array (0..7) of Boolean;
29     type State_Mask  is array (State) of Boolean;
30     type Mode_Mask   is array (Mode) of Boolean;
31     type Program_Status_Word is
32       record
33         System_Mask      : Byte_Mask;
34         Protection_Key    : Integer range 0 .. 3;
35         Machine_State     : State_Mask;
36         Interrupt_Cause   : Interruption_Code;
37         Ilc                : Integer range 0 .. 3;
38         Cc                 : Integer range 0 .. 3;
39         Program_Mask      : Mode_Mask;
40         Inst_Address       : Address;
41       end record;
42     for Program_Status_Word use
43       record
44         System_Mask      at 0*Word range 0 .. 7;
45         Protection_Key   at 0*Word range 10 .. 11; -- bits 8,9 unused
46         Machine_State     at 0*Word range 12 .. 15;
47         Interrupt_Cause   at 0*Word range 16 .. 31;
48         Ilc                at 1*Word range 0 .. 1; -- second word
49         Cc                 at 1*Word range 2 .. 3;
50         Program_Mask      at 1*Word range 4 .. 7;
51         Inst_Address       at 1*Word range 8 .. 31;
52       end record;
53     for Program_Status_Word'Size use 8*System.Storage_Unit;
54     for Program_Status_Word'Alignment use 8;
```

NOTES

31 12 *Note on the example:* The `record_representation_clause` defines the record layout. The `Size` clause guarantees that (at least) eight storage elements are used for objects of the type. The `Alignment` clause guarantees that aliased, imported, or exported objects of the type will have addresses divisible by eight.

13.5.2 Storage Place Attributes

Static Semantics

1 For a component C of a composite, non-array object R, the *storage place attributes* are defined:

2/2 R.C'Position If the nondefault bit ordering applies to the composite type, and if a `component_clause` specifies the placement of C, denotes the value given for the `position` of the `component_clause`; otherwise, denotes the same value as R.C'Address – R'Address. The value of this attribute is of the type `universal_integer`.

3/2 R.C'First_Bit

If the nondefault bit ordering applies to the composite type, and if a `component_clause` specifies the placement of C, denotes the value given for the `first_bit` of the `component_clause`; otherwise, denotes the offset, from the start of the first of the storage elements occupied by C, of the first bit occupied by C. This offset is measured in bits. The

first bit of a storage element is numbered zero. The value of this attribute is of the type *universal_integer*.

R.CLast_Bit

If the nondefault bit ordering applies to the composite type, and if a *component_clause* specifies the placement of C, denotes the value given for the *last_bit* of the *component_clause*; otherwise, denotes the offset, from the start of the first of the storage elements occupied by C, of the last bit occupied by C. This offset is measured in bits. The value of this attribute is of the type *universal_integer*.

Implementation Advice

If a component is represented using some form of pointer (such as an offset) to the actual data of the component, and this data is contiguous with the rest of the object, then the storage place attributes should reflect the place of the actual data, not the pointer. If a component is allocated discontiguously from the rest of the object, then a warning should be generated upon reference to one of its storage place attributes.

4/2

5

13.5.3 Bit Ordering

The *Bit_Order* attribute specifies the interpretation of the storage place attributes.

1

Static Semantics

A bit ordering is a method of interpreting the meaning of the storage place attributes. *High_Order_First* (known in the vernacular as “big endian”) means that the first bit of a storage element (bit 0) is the most significant bit (interpreting the sequence of bits that represent a component as an unsigned integer value). *Low_Order_First* (known in the vernacular as “little endian”) means the opposite: the first bit is the least significant.

2

For every specific record subtype S, the following attribute is defined:

3

S'Bit_Order Denotes the bit ordering for the type of S. The value of this attribute is of type *System.Bit_Order*. *Bit_Order* may be specified for specific record types via an *attribute_definition_clause*; the expression of such a clause shall be static.

4

If *Word_Size* = *Storage_Unit*, the default bit ordering is implementation defined. If *Word_Size* > *Storage_Unit*, the default bit ordering is the same as the ordering of storage elements in a word, when interpreted as an integer.

5

The storage place attributes of a component of a type are interpreted according to the bit ordering of the type.

6

Implementation Advice

The recommended level of support for the nondefault bit ordering is:

7

- The implementation should support the nondefault bit ordering in addition to the default bit ordering.

8/2

NOTES

13 *Bit_Order* clauses make it possible to write *record_representation_clauses* that can be ported between machines having different bit ordering. They do not guarantee transparent exchange of data between such machines.

9/2

13.6 Change of Representation

- 1 A type_conversion (see 4.6) can be used to convert between two different representations of the same array or record. To convert an array from one representation to another, two array types need to be declared with matching component subtypes, and convertible index types. If one type has packing specified and the other does not, then explicit conversion can be used to pack or unpack an array.
- 2 To convert a record from one representation to another, two record types with a common ancestor type need to be declared, with no inherited subprograms. Distinct representations can then be specified for the record types, and explicit conversion between the types can be used to effect a change in representation.

Examples

```

3   Example of change of representation:
4     -- Packed_Descriptor and Descriptor are two different types
4     -- with identical characteristics, apart from their
4     -- representation
5
6     type Descriptor is
7       record
8         -- components of a descriptor
9       end record;
10
11    type Packed_Descriptor is new Descriptor;
12    for Packed_Descriptor use
13      record
14        -- component clauses for some or for all components
15      end record;
16
17    -- Change of representation can now be accomplished by explicit type conversions:
18    D : Descriptor;
19    P : Packed_Descriptor;
20
21    P := Packed_Descriptor(D);    -- pack D
22    D := Descriptor(P);          -- unpack P

```

13.7 The Package System

For each implementation there is a library package called System which includes the definitions of certain configuration-dependent characteristics.

Static Semantics

The following language-defined library package exists:

```

package System is
  pragma Pure(System);
  type Name is implementation-defined-enumeration-type;
  System_Name : constant Name := implementation-defined;
  -- System-Dependent Named Numbers:
  Min_Int           : constant := root_integer'First;
  Max_Int           : constant := root_integer>Last;
  Max_Binary_Modulus : constant := implementation-defined;
  Max_Nonbinary_Modulus : constant := implementation-defined;
  Max_Base_Digits   : constant := root_real'Digits;
  Max_Digits         : constant := implementation-defined;
  Max_Mantissa       : constant := implementation-defined;
  Fine_Delta          : constant := implementation-defined;
  Tick               : constant := implementation-defined;
  -- Storage-related Declarations:
  type Address is implementation-defined;
  Null_Address : constant Address;
  Storage_Unit    : constant := implementation-defined;
  Word_Size        : constant := implementation-defined * Storage_Unit;
  Memory_Size      : constant := implementation-defined;
  -- Address Comparison:
  function "<" (Left, Right : Address) return Boolean;
  function "<=" (Left, Right : Address) return Boolean;
  function ">" (Left, Right : Address) return Boolean;
  function ">=" (Left, Right : Address) return Boolean;
  function "=" (Left, Right : Address) return Boolean;
  -- function "/=" (Left, Right : Address) return Boolean;
  -- "/=" is implicitly defined
  pragma Convention(Intrinsic, "<");
  ... -- and so on for all language-defined subprograms in this package
  -- Other System-Dependent Declarations:
  type Bit_Order is (High_Order_First, Low_Order_First);
  Default_Bit_Order : constant Bit_Order := implementation-defined;
  -- Priority-related declarations (see D.1):
  subtype Any_Priority is Integer range implementation-defined;
  subtype Priority is Any_Priority range Any_Priority'First .. implementation-defined;
  subtype Interrupt_Priority is Any_Priority range Priority'Last+1 .. Any_Priority'Last;
  Default_Priority : constant Priority := (Priority'First + Priority'Last)/2;
  private
    ... -- not specified by the language
  end System;
```

Name is an enumeration subtype. Values of type Name are the names of alternative machine configurations handled by the implementation. System_Name represents the current machine configuration.

- 20 The named numbers `Fine_Delta` and `Tick` are of the type `universal_real`; the others are of the type `universal_integer`.
- 21 The meanings of the named numbers are:
 - 22 `Min_Int` The smallest (most negative) value allowed for the expressions of a `signed_integer_type_definition`.
 - 23 `Max_Int` The largest (most positive) value allowed for the expressions of a `signed_integer_type_definition`.
 - 24 `Max_Binary_Modulus`
 - A power of two such that it, and all lesser positive powers of two, are allowed as the modulus of a `modular_type_definition`.
 - 25 `Max_Nonbinary_Modulus`
 - A value such that it, and all lesser positive integers, are allowed as the modulus of a `modular_type_definition`.
 - 26 `Max_Base_Digits`
 - The largest value allowed for the requested decimal precision in a `floating_point_definition`.
 - 27 `Max_Digits` The largest value allowed for the requested decimal precision in a `floating_point_definition` that has no `real_range_specification`. `Max_Digits` is less than or equal to `Max_Base_Digits`.
 - 28 `Max_Mantissa`
 - The largest possible number of binary digits in the mantissa of machine numbers of a user-defined ordinary fixed point type. (The mantissa is defined in Annex G.)
 - 29 `Fine_Delta` The smallest delta allowed in an `ordinary_fixed_point_definition` that has the `real_range_specification range -1.0 .. 1.0`.
 - 30 `Tick` A period in seconds approximating the real time interval during which the value of `Calendar.Clock` remains constant.
 - 31 `Storage_Unit`
 - The number of bits per storage element.
 - 32 `Word_Size` The number of bits per word.
 - 33 `Memory_Size` An implementation-defined value that is intended to reflect the memory size of the configuration in storage elements.
 - 34/2 Address is a definite, nonlimited type with preelaborable initialization (see 10.2.1). Address represents machine addresses capable of addressing individual storage elements. `Null_Address` is an address that is distinct from the address of any object or program unit.
 - 35/2 `Default_Bit_Order` shall be a static constant. See 13.5.3 for an explanation of `Bit_Order` and `Default_Bit_Order`.

Implementation Permissions

- 36/2 An implementation may add additional implementation-defined declarations to package `System` and its children. However, it is usually better for the implementation to provide additional functionality via implementation-defined children of `System`.

Implementation Advice

- 37 Address should be a private type.

NOTES

14 There are also some language-defined child packages of System defined elsewhere.

38

13.7.1 The Package System.Storage_Elements

Static Semantics

The following language-defined library package exists:

```

package System.Storage_Elements is
  pragma Pure(Storage_Elements);

  type Storage_Offset is range implementation-defined;
  subtype Storage_Count is Storage_Offset range 0..Storage_Offset'Last;
  type Storage_Element is mod implementation-defined;
  for Storage_Element'Size use Storage_Unit;
  type Storage_Array is array
    (Storage_Offset range <>) of aliased Storage_Element;
  for Storage_Array'Component_Size use Storage_Unit;

  -- Address Arithmetic:
  function "+"(Left : Address; Right : Storage_Offset)
    return Address;
  function "+"(Left : Storage_Offset; Right : Address)
    return Address;
  function "-"(Left : Address; Right : Storage_Offset)
    return Address;
  function "-"(Left, Right : Address)
    return Storage_Offset;
  function "mod"(Left : Address; Right : Storage_Offset)
    return Storage_Offset;
  -- Conversion to/from integers:
  type Integer_Address is implementation-defined;
  function To_Address(Value : Integer_Address) return Address;
  function To_Integer(Value : Address) return Integer_Address;
  pragma Convention(Intrinsic, "+");
  --...and so on for all language-defined subprograms declared in this package.
end System.Storage_Elements;

```

Storage_Element represents a storage element. Storage_Offset represents an offset in storage elements. Storage_Count represents a number of storage elements. Storage_Array represents a contiguous sequence of storage elements.

Integer_Address is a (signed or modular) integer subtype. To_Address and To_Integer convert back and forth between this type and Address.

Implementation Requirements

Storage_Offset'Last shall be greater than or equal to Integer'Last or the largest possible storage offset, whichever is smaller. Storage_Offset'First shall be <= (-Storage_Offset'Last).

Implementation Permissions

This paragraph was deleted.

15/2

Implementation Advice

Operations in System and its children should reflect the target environment semantics as closely as is reasonable. For example, on most machines, it makes sense for address arithmetic to “wrap around.” Operations that do not make sense should raise Program_Error.

13.7.2 The Package System.Address_To_Access_Conversions

Static Semantics

- 1 The following language-defined generic library package exists:

```

2   generic
3     type Object(<>) is limited private;
4   package System.Address_To_Access_Conversions is
5     pragma Preelaborate(Address_To_Access_Conversions);
6     type Object_Pointer is access all Object;
7     function To_Pointer(Value : Address) return Object_Pointer;
8     function To_Address(Value : Object_Pointer) return Address;
9     pragma Convention(Intrinsic, To_Pointer);
10    pragma Convention(Intrinsic, To_Address);
11   end System.Address_To_Access_Conversions;

```

- 5/2 The To_Pointer and To_Address subprograms convert back and forth between values of types Object_Pointer and Address. To_Pointer(X'Address) is equal to X'Unchecked_Access for any X that allows Unchecked_Access. To_Pointer(Null_Address) returns null. For other addresses, the behavior is unspecified. To_Address(null) returns Null_Address. To_Address(Y), where Y /= null, returns Y.all'Address.

Implementation Permissions

- 6 An implementation may place restrictions on instantiations of Address_To_Access_Conversions.

13.8 Machine Code Insertions

- 1 A machine code insertion can be achieved by a call to a subprogram whose sequence_of_statements contains code_statements.

Syntax

2 code_statement ::= qualified_expression;

3 A code_statement is only allowed in the handled_sequence_of_statements of a subprogram_body. If a subprogram_body contains any code_statements, then within this subprogram_body the only allowed form of statement is a code_statement (labeled or not), the only allowed declarative_items are use_clauses, and no exception_handler is allowed (comments and pragmas are allowed as usual).

Name Resolution Rules

- 4 The qualified_expression is expected to be of any type.

Legality Rules

- 5 The qualified_expression shall be of a type declared in package System.Machine_Code.
- 6 A code_statement shall appear only within the scope of a with_clause that mentions package System.Machine_Code.

Static Semantics

- 7 The contents of the library package System.Machine_Code (if provided) are implementation defined. The meaning of code_statements is implementation defined. Typically, each qualified_expression represents a machine instruction or assembly directive.

Implementation Permissions

An implementation may place restrictions on `code_statements`. An implementation is not required to provide package `System.Machine_Code`.

NOTES

15 An implementation may provide implementation-defined pragmas specifying register conventions and calling conventions.

16 Machine code functions are exempt from the rule that a return statement is required. In fact, return statements are forbidden, since only `code_statements` are allowed.

17 Intrinsic subprograms (see 6.3.1, “Conformance Rules”) can also be used to achieve machine code insertions. Interface to assembly language can be achieved using the features in Annex B, “Interface to Other Languages”.

Examples

Example of a code statement:

```
M : Mask;
procedure Set_Mask; pragma Inline(Set_Mask);
procedure Set_Mask is
  use System.Machine_Code; -- assume "with System.Machine_Code;" appears somewhere above
begin
  SI_Format'(Code => SSM, B => M'Base_Reg, D => M'Disp);
  -- Base_Reg and Disp are implementation-defined attributes
end Set_Mask;
```

13.9 Unchecked Type Conversions

An unchecked type conversion can be achieved by a call to an instance of the generic function `Unchecked_Conversion`.

Static Semantics

The following language-defined generic library function exists:

```
generic
  type Source(<>) is limited private;
  type Target(<>) is limited private;
function Ada.Unchecked_Conversion(S : Source) return Target;
pragma Convention(Intrinsic, Ada.Unchecked_Conversion);
pragma Pure(Ada.Unchecked_Conversion);
```

Dynamic Semantics

The size of the formal parameter `S` in an instance of `Unchecked_Conversion` is that of its subtype. This is the actual subtype passed to `Source`, except when the actual is an unconstrained composite subtype, in which case the subtype is constrained by the bounds or discriminants of the value of the actual expression passed to `S`.

If all of the following are true, the effect of an unchecked conversion is to return the value of an object of the target subtype whose representation is the same as that of the source object `S`:

- `S'Size = Target'Size`.
- `S'Alignment = Target'Alignment`.
- The target subtype is not an unconstrained composite subtype.
- `S` and the target subtype both have a contiguous representation.
- The representation of `S` is a representation of an object of the target subtype.

- 11/2 Otherwise, if the result type is scalar, the result of the function is implementation defined, and can have an invalid representation (see 13.9.1). If the result type is nonscalar, the effect is implementation defined; in particular, the result can be abnormal (see 13.9.1).

Implementation Permissions

- 12 An implementation may return the result of an unchecked conversion by reference, if the Source type is not a by-copy type. In this case, the result of the unchecked conversion represents simply a different (read-only) view of the operand of the conversion.
- 13 An implementation may place restrictions on `Unchecked_Conversion`.

Implementation Advice

- 14/2 Since the `Size` of an array object generally does not include its bounds, the bounds should not be part of the converted data.
- 15 The implementation should not generate unnecessary run-time checks to ensure that the representation of S is a representation of the target type. It should take advantage of the permission to return by reference when possible. Restrictions on unchecked conversions should be avoided unless required by the target environment.
- 16 The recommended level of support for unchecked conversions is:
- 17 • Unchecked conversions should be supported and should be reversible in the cases where this clause defines the result. To enable meaningful use of unchecked conversion, a contiguous representation should be used for elementary subtypes, for statically constrained array subtypes whose component subtype is one of the subtypes described in this paragraph, and for record subtypes without discriminants whose component subtypes are described in this paragraph.

13.9.1 Data Validity

- 1 Certain actions that can potentially lead to erroneous execution are not directly erroneous, but instead can cause objects to become *abnormal*. Subsequent uses of abnormal objects can be erroneous.
- 2 A scalar object can have an *invalid representation*, which means that the object's representation does not represent any value of the object's subtype. The primary cause of invalid representations is uninitialized variables.
- 3 Abnormal objects and invalid representations are explained in this subclause.

Dynamic Semantics

- 4 When an object is first created, and any explicit or default initializations have been performed, the object and all of its parts are in the *normal* state. Subsequent operations generally leave them normal. However, an object or part of an object can become *abnormal* in the following ways:
- 5 • An assignment to the object is disrupted due to an abort (see 9.8) or due to the failure of a language-defined check (see 11.6).
- 6/2 • The object is not scalar, and is passed to an `in out` or `out` parameter of an imported procedure, the `Read` procedure of an instance of `Sequential_IO`, `Direct_IO`, or `Storage_IO`, or the stream attribute `T'Read`, if after return from the procedure the representation of the parameter does not represent a value of the parameter's subtype.
- 6.1/2 • The object is the return object of a function call of a nonscalar type, and the function is an imported function, an instance of `Unchecked_Conversion`, or the stream attribute `T'Input`, if after

return from the function the representation of the return object does not represent a value of the function's subtype.

For an imported object, it is the programmer's responsibility to ensure that the object remains in a normal state. 6.2/2

Whether or not an object actually becomes abnormal in these cases is not specified. An abnormal object becomes normal again upon successful completion of an assignment to the object as a whole. 7

Erroneous Execution

It is erroneous to evaluate a primary that is a **name** denoting an abnormal object, or to evaluate a prefix that denotes an abnormal object. 8

Bounded (Run-Time) Errors

If the representation of a scalar object does not represent a value of the object's subtype (perhaps because the object was not initialized), the object is said to have an *invalid representation*. It is a bounded error to evaluate the value of such an object. If the error is detected, either **Constraint_Error** or **Program_Error** is raised. Otherwise, execution continues using the invalid representation. The rules of the language outside this subclause assume that all objects have valid representations. The semantics of operations on invalid representations are as follows:

- If the representation of the object represents a value of the object's type, the value of the type is used. 10
- If the representation of the object does not represent a value of the object's type, the semantics of operations on such representations is implementation-defined, but does not by itself lead to erroneous or unpredictable execution, or to other objects becoming abnormal. 11

Erroneous Execution

A call to an imported function or an instance of **Unchecked_Conversion** is erroneous if the result is scalar, the result object has an invalid representation, and the result is used other than as the **expression** of an **assignment_statement** or an **object_declaration**, or as the **prefix** of a **Valid** attribute. If such a result object is used as the source of an assignment, and the assigned value is an invalid representation for the target of the assignment, then any use of the target object prior to a further assignment to the target object, other than as the **prefix** of a **Valid** attribute reference, is erroneous. 12/2

The dereference of an access value is erroneous if it does not designate an object of an appropriate type or a subprogram with an appropriate profile, if it designates a nonexistent object, or if it is an access-to-variable value that designates a constant object. Such an access value can exist, for example, because of **Unchecked_Deallocation**, **Unchecked_Access**, or **Unchecked_Conversion**. 13

NOTES

18 Objects can become abnormal due to other kinds of actions that directly update the object's representation; such actions are generally considered directly erroneous, however. 14

13.9.2 The Valid Attribute

The **Valid** attribute can be used to check the validity of data produced by unchecked conversion, input, interface to foreign languages, and the like. 1

Static Semantics

For a prefix **X** that denotes a scalar object (after any implicit dereference), the following attribute is defined: 2

- 3 X'Valid Yields True if and only if the object denoted by X is normal and has a valid representation.
The value of this attribute is of the predefined type Boolean.

NOTES

4 19 Invalid data can be created in the following cases (not counting erroneous or unpredictable execution):

- 5 • an uninitialized scalar object,
- 6 • the result of an unchecked conversion,
- 7 • input,
- 8 • interface to another language (including machine code),
- 9 • aborting an assignment,
- 10 • disrupting an assignment due to the failure of a language-defined check (see 11.6), and
- 11 • use of an object whose Address has been specified.

12 20 X'Valid is not considered to be a read of X; hence, it is not an error to check the validity of invalid data.

13/2 21 The Valid attribute may be used to check the result of calling an instance of `Unchecked_Conversion` (or any other operation that can return invalid values). However, an exception handler should also be provided because implementations are permitted to raise `Constraint_Error` or `Program_Error` if they detect the use of an invalid representation (see 13.9.1).

13.10 Unchecked Access Value Creation

- 1 The attribute `Unchecked_Access` is used to create access values in an unsafe manner — the programmer is responsible for preventing “dangling references.”

Static Semantics

- 2 The following attribute is defined for a prefix X that denotes an aliased view of an object:

3 X'Unchecked_Access

All rules and semantics that apply to `X'Access` (see 3.10.2) apply also to `X'Unchecked_Access`, except that, for the purposes of accessibility rules and checks, it is as if X were declared immediately within a library package.

NOTES

4 22 This attribute is provided to support the situation where a local object is to be inserted into a global linked data structure, when the programmer knows that it will always be removed from the data structure prior to exiting the object's scope. The `Access` attribute would be illegal in this case (see 3.10.2, “Operations of Access Types”).

5 23 There is no `Unchecked_Access` attribute for subprograms.

13.11 Storage Management

- 1 Each access-to-object type has an associated storage pool. The storage allocated by an allocator comes from the pool; instances of `Unchecked_Deallocation` return storage to the pool. Several access types can share the same pool.

2/2 A storage pool is a variable of a type in the class rooted at `Root_Storage_Pool`, which is an abstract limited controlled type. By default, the implementation chooses a *standard storage pool* for each access-to-object type. The user may define new pool types, and may override the choice of pool for an access-to-object type by specifying `Storage_Pool` for the type.

Legality Rules

- 3 If `Storage_Pool` is specified for a given access type, `Storage_Size` shall not be specified for it.

Static Semantics

The following language-defined library package exists:

```

with Ada.Finalization;
with System.Storage_Elements;
package System.Storage_Pools is
  pragma Preelaborate(System.Storage_Pools);
  type Root_Storage_Pool is
    abstract new Ada.Finalization.Limited_Controlled with private;
  pragma Preelaborate_Initialization(Root_Storage_Pool);
  procedure Allocate(
    Pool : in out Root_Storage_Pool;
    Storage_Address : out Address;
    Size_In_Storage_Elements : in Storage_Elements.Storage_Count;
    Alignment : in Storage_Elements.Storage_Count) is abstract;
  procedure Deallocate(
    Pool : in out Root_Storage_Pool;
    Storage_Address : in Address;
    Size_In_Storage_Elements : in Storage_Elements.Storage_Count;
    Alignment : in Storage_Elements.Storage_Count) is abstract;
  function Storage_Size(Pool : Root_Storage_Pool)
    return Storage_Elements.Storage_Count is abstract;
private
  ... -- not specified by the language
end System.Storage_Pools;

```

A *storage pool type* (or *pool type*) is a descendant of `Root_Storage_Pool`. The *elements* of a storage pool are the objects allocated in the pool by allocators.

For every access-to-object subtype S, the following representation attributes are defined:

`S'Storage_Pool`

Denotes the storage pool of the type of S. The type of this attribute is `Root_Storage_Pool'Class`.

`S'Storage_Size`

Yields the result of calling `Storage_Size(S'Storage_Pool)`, which is intended to be a measure of the number of storage elements reserved for the pool. The type of this attribute is *universal_integer*.

`Storage_Size` or `Storage_Pool` may be specified for a non-derived access-to-object type via an `attribute_definition_clause`; the name in a `Storage_Pool` clause shall denote a variable.

An allocator of type T allocates storage from T's storage pool. If the storage pool is a user-defined object, then the storage is allocated by calling `Allocate`, passing `T'Storage_Pool` as the `Pool` parameter. The `Size_In_Storage_Elements` parameter indicates the number of storage elements to be allocated, and is no more than `D'Max_Size_In_Storage_Elements`, where D is the designated subtype. The `Alignment` parameter is `D'Alignment`. The result returned in the `Storage_Address` parameter is used by the allocator as the address of the allocated storage, which is a contiguous block of memory of `Size_In_Storage_Elements` storage elements. Any exception propagated by `Allocate` is propagated by the allocator.

If `Storage_Pool` is not specified for a type defined by an `access_to_object_definition`, then the implementation chooses a standard storage pool for it in an implementation-defined manner. In this case, the exception `Storage_Error` is raised by an allocator if there is not enough storage. It is implementation defined whether or not the implementation provides user-accessible names for the standard pool type(s).

If `Storage_Size` is specified for an access type, then the `Storage_Size` of this pool is at least that requested, and the storage for the pool is reclaimed when the master containing the declaration of the access type is

left. If the implementation cannot satisfy the request, Storage_Error is raised at the point of the attribute_definition_clause. If neither Storage_Pool nor Storage_Size are specified, then the meaning of Storage_Size is implementation defined.

- 19 If Storage_Pool is specified for an access type, then the specified pool is used.
- 20 The effect of calling Allocate and Deallocate for a standard storage pool directly (rather than implicitly via an allocator or an instance of Unchecked_Deallocation) is unspecified.

Erroneous Execution

- 21 If Storage_Pool is specified for an access type, then if Allocate can satisfy the request, it should allocate a contiguous block of memory, and return the address of the first storage element in Storage_Address. The block should contain Size_In_Storage_Elements storage elements, and should be aligned according to Alignment. The allocated storage should not be used for any other purpose while the pool element remains in existence. If the request cannot be satisfied, then Allocate should propagate an exception (such as Storage_Error). If Allocate behaves in any other manner, then the program execution is erroneous.

Documentation Requirements

- 22 An implementation shall document the set of values that a user-defined Allocate procedure needs to accept for the Alignment parameter. An implementation shall document how the standard storage pool is chosen, and how storage is allocated by standard storage pools.

Implementation Advice

- 23 An implementation should document any cases in which it dynamically allocates heap storage for a purpose other than the evaluation of an allocator.
- 24 A default (implementation-provided) storage pool for an access-to-constant type should not have overhead to support deallocation of individual objects.
- 25/2 The storage pool used for an allocator of an anonymous access type should be determined as follows:
 - If the allocator is defining a coextension (see 3.10.2) of an object being created by an outer allocator, then the storage pool used for the outer allocator should also be used for the coextension;
 - For other access discriminants and access parameters, the storage pool should be created at the point of the allocator, and be reclaimed when the allocated object becomes inaccessible;
 - Otherwise, a default storage pool should be created at the point where the anonymous access type is elaborated; such a storage pool need not support deallocation of individual objects.

NOTES

- 26 24 A user-defined storage pool type can be obtained by extending the Root_Storage_Pool type, and overriding the primitive subprograms Allocate, Deallocate, and Storage_Size. A user-defined storage pool can then be obtained by declaring an object of the type extension. The user can override Initialize and Finalize if there is any need for non-trivial initialization and finalization for a user-defined pool type. For example, Finalize might reclaim blocks of storage that are allocated separately from the pool object itself.
- 27 25 The writer of the user-defined allocation and deallocation procedures, and users of allocators for the associated access type, are responsible for dealing with any interactions with tasking. In particular:
 - If the allocators are used in different tasks, they require mutual exclusion.
 - If they are used inside protected objects, they cannot block.
 - If they are used by interrupt handlers (see C.3, “Interrupt Support”), the mutual exclusion mechanism has to work properly in that context.

26 The primitives Allocate, Deallocate, and Storage_Size are declared as abstract (see 3.9.3), and therefore they have to
be overridden when a new (non-abstract) storage pool type is declared.

Examples

To associate an access type with a storage pool object, the user first declares a pool object of some type
derived from Root_Storage_Pool. Then, the user defines its Storage_Pool attribute, as follows:

```
Pool_Object : Some_Storage_Pool_Type;
type T is access Designated;
for T'Storage_Pool use Pool_Object;
```

Another access type may be added to an existing storage pool, via:

```
for T2'Storage_Pool use T'Storage_Pool;
```

The semantics of this is implementation defined for a standard storage pool.

As usual, a derivative of Root_Storage_Pool may define additional operations. For example, presuming
that Mark_Release_Pool_Type has two additional operations, Mark and Release, the following is a
possible use:

```
type Mark_Release_Pool_Type
  (Pool_Size : Storage_Elements.Storage_Count;
   Block_Size : Storage_Elements.Storage_Count)
  is new Root_Storage_Pool with private;
...
MR_Pool : Mark_Release_Pool_Type (Pool_Size => 2000,
                                   Block_Size => 100);
type Acc is access ...;
for Acc'Storage_Pool use MR_Pool;
...
Mark(MR_Pool);
... -- Allocate objects using "new Designated(...)".
Release(MR_Pool); -- Reclaim the storage.
```

13.11.1 The Max_Size_In_Storage_Elements Attribute

The Max_Size_In_Storage_Elements attribute is useful in writing user-defined pool types.

Static Semantics

For every subtype S, the following attribute is defined:

S'Max_Size_In_Storage_Elements

Denotes the maximum value for Size_In_Storage_Elements that could be requested by the
implementation via Allocate for an access type whose designated subtype is S. For a type
with access discriminants, if the implementation allocates space for a coextension in the
same pool as that of the object having the access discriminant, then this accounts for any
calls on Allocate that could be performed to provide space for such coextensions. The value
of this attribute is of type *universal_integer*.

13.11.2 Unchecked Storage Deallocation

- 1 Unchecked storage deallocation of an object designated by a value of an access type is achieved by a call to an instance of the generic procedure `Unchecked_Deallocation`.

Static Semantics

- 2 The following language-defined generic library procedure exists:

```
3   generic
4     type Object(<>) is limited private;
5     type Name    is access Object;
6     procedure Ada.Unchecked_Deallocation(X : in out Name);
7     pragma Convention(Intrinsic, Ada.Unchecked_Deallocation);
8     pragma Preelaborate(Ada.Unchecked_Deallocation);
```

Dynamic Semantics

- 4 Given an instance of `Unchecked_Deallocation` declared as follows:

```
5   procedure Free is
6     new Ada.Unchecked_Deallocation(
7       object_subtype_name, access_to_variable_subtype_name);
```

- 6 Procedure `Free` has the following effect:

- 7 1. After executing `Free(X)`, the value of `X` is **null**.
- 8 2. `Free(X)`, when `X` is already equal to **null**, has no effect.
- 9/2 3. `Free(X)`, when `X` is not equal to **null** first performs finalization of the object designated by `X` (and any coextensions of the object — see 3.10.2), as described in 7.6.1. It then deallocates the storage occupied by the object designated by `X` (and any coextensions). If the storage pool is a user-defined object, then the storage is deallocated by calling `Deallocate`, passing `access_to_variable_subtype_name'Storage_Pool` as the `Pool` parameter. `Storage_Address` is the value returned in the `Storage_Address` parameter of the corresponding `Allocate` call. `Size_In_Storage_Elements` and `Alignment` are the same values passed to the corresponding `Allocate` call. There is one exception: if the object being freed contains tasks, the object might not be deallocated.

- 10/2 After `Free(X)`, the object designated by `X`, and any subcomponents (and coextensions) thereof, no longer exist; their storage can be reused for other purposes.

Bounded (Run-Time) Errors

- 11 It is a bounded error to free a discriminated, unterminated task object. The possible consequences are:

- 12 • No exception is raised.
- 13 • `Program_Error` or `Tasking_Error` is raised at the point of the deallocation.
- 14 • `Program_Error` or `Tasking_Error` is raised in the task the next time it references any of the discriminants.

- 15 In the first two cases, the storage for the discriminants (and for any enclosing object if it is designated by an access discriminant of the task) is not reclaimed prior to task termination.

Erroneous Execution

- 16 Evaluating a name that denotes a nonexistent object is erroneous. The execution of a call to an instance of `Unchecked_Deallocation` is erroneous if the object was created other than by an allocator for an access type whose pool is `Name'Storage_Pool`.

Implementation Advice

For a standard storage pool, Free should actually reclaim the storage.

17

NOTES

27 The rules here that refer to Free apply to any instance of Unchecked_Deallocation.

18

28 Unchecked_Deallocation cannot be instantiated for an access-to-constant type. This is implied by the rules of 12.5.4.

19

13.11.3 Pragma Controlled

Pragma Controlled is used to prevent any automatic reclamation of storage (garbage collection) for the objects created by allocators of a given access type.

1

Syntax

The form of a **pragma** Controlled is as follows:

2

pragma Controlled(*first_subtype_local_name*);

3

Legality Rules

The *first_subtype_local_name* of a **pragma** Controlled shall denote a non-derived access subtype.

4

Static Semantics

A **pragma** Controlled is a representation pragma that specifies the *controlled* aspect of representation.

5

Garbage collection is a process that automatically reclaims storage, or moves objects to a different address, while the objects still exist.

6

If a **pragma** Controlled is specified for an access type with a standard storage pool, then garbage collection is not performed for objects in that pool.

7

Implementation Permissions

An implementation need not support garbage collection, in which case, a **pragma** Controlled has no effect.

8

13.12 Pragma Restrictions

A **pragma** Restrictions expresses the user's intent to abide by certain restrictions. This may facilitate the construction of simpler run-time environments.

1

Syntax

The form of a **pragma** Restrictions is as follows:

2

pragma Restrictions(*restriction* {, *restriction*});

3

restriction ::= *restriction_identifier*

4/2

| *restriction_parameter_identifier* => *restriction_parameter_argument*

restriction_parameter_argument ::= *name* | *expression*

4.1/2

Name Resolution Rules

Unless otherwise specified for a particular restriction, the **expression** is expected to be of any integer type.

5

Legality Rules

Unless otherwise specified for a particular restriction, the **expression** shall be static, and its value shall be nonnegative.

6

Static Semantics

- 7/2 The set of restrictions is implementation defined.

Post-Compilation Rules

- 8 A **pragma Restrictions** is a configuration pragma; unless otherwise specified for a particular restriction, a partition shall obey the restriction if a **pragma Restrictions** applies to any compilation unit included in the partition.
- 8.1/1 For the purpose of checking whether a partition contains constructs that violate any restriction (unless specified otherwise for a particular restriction):
- 8.2/1 • Generic instances are logically expanded at the point of instantiation;
 - 8.3/1 • If an object of a type is declared or allocated and not explicitly initialized, then all expressions appearing in the definition for the type and any of its ancestors are presumed to be used;
 - 8.4/1 • A **default_expression** for a formal parameter or a generic formal object is considered to be used if and only if the corresponding actual parameter is not provided in a given call or instantiation.

Implementation Permissions

- 9 An implementation may place limitations on the values of the **expression** that are supported, and limitations on the supported combinations of restrictions. The consequences of violating such limitations are implementation defined.
- 9.1/1 An implementation is permitted to omit restriction checks for code that is recognized at compile time to be unreachable and for which no code is generated.
- 9.2/1 Whenever enforcement of a restriction is not required prior to execution, an implementation may nevertheless enforce the restriction prior to execution of a partition to which the restriction applies, provided that every execution of the partition would violate the restriction.

NOTES

- 10/2 29 Restrictions intended to facilitate the construction of efficient tasking run-time systems are defined in D.7. Restrictions intended for use when constructing high integrity systems are defined in H.4.
- 11 30 An implementation has to enforce the restrictions in cases where enforcement is required, even if it chooses not to take advantage of the restrictions in terms of efficiency.

13.12.1 Language-Defined Restrictions

Static Semantics

- 1/2 The following **restriction_identifiers** are language-defined (additional restrictions are defined in the Specialized Needs Annexes):
- 2/2 **No_Implementation_Attributes**
There are no implementation-defined attributes. This restriction applies only to the current compilation or environment, not the entire partition.
- 3/2 **No_Implementation_Pragmas**
There are no implementation-defined pragmas or pragma arguments. This restriction applies only to the current compilation or environment, not the entire partition.
- 4/2 **No_Obsolescent_Features**
There is no use of language features defined in Annex J. It is implementation-defined if uses of the renamings of J.1 are detected by this restriction. This restriction applies only to the current compilation or environment, not the entire partition.

The following *restriction_parameter_identifier* is language defined:

No_Dependence

Specifies a library unit on which there are no semantic dependences.

Legality Rules

The *restriction_parameter_argument* of a No_Dependence restriction shall be a name; the name shall have the form of a full expanded name of a library unit, but need not denote a unit present in the environment.

Post-Compilation Rules

No compilation unit included in the partition shall depend semantically on the library unit identified by the name.

13.13 Streams

A *stream* is a sequence of elements comprising values from possibly different types and allowing sequential access to these values. A *stream type* is a type in the class whose root type is Streams.Root_Stream_Type. A stream type may be implemented in various ways, such as an external sequential file, an internal buffer, or a network channel.

13.13.1 The Package Streams

Static Semantics

The abstract type Root_Stream_Type is the root type of the class of stream types. The types in this class represent different kinds of streams. A new stream type is defined by extending the root type (or some other stream type), overriding the Read and Write operations, and optionally defining additional primitive subprograms, according to the requirements of the particular kind of stream. The predefined stream-oriented attributes like T'Read and T'Write make dispatching calls on the Read and Write procedures of the Root_Stream_Type. (User-defined T'Read and T'Write attributes can also make such calls, or can call the Read and Write attributes of other types.)

```

package Ada.Streams is
  pragma Pure(Streams);
  type Root_Stream_Type is abstract tagged limited private;
  pragma Preelaborable_Initialization(Root_Stream_Type);
  type Stream_Element is mod implementation-defined;
  type Stream_Element_Offset is range implementation-defined;
  subtype Stream_Element_Count is
    Stream_Element_Offset range 0 .. Stream_Element_Offset'Last;
  type Stream_Element_Array is
    array(Stream_Element_Offset range <>) of aliased Stream_Element;
  procedure Read(
    Stream : in out Root_Stream_Type;
    Item   : out Stream_Element_Array;
    Last   : out Stream_Element_Offset) is abstract;
  procedure Write(
    Stream : in out Root_Stream_Type;
    Item   : in Stream_Element_Array) is abstract;
private
  ... -- not specified by the language
end Ada.Streams;

```

- 8/2 The Read operation transfers stream elements from the specified stream to fill the array Item. Elements are transferred until Item'Length elements have been transferred, or until the end of the stream is reached. If any elements are transferred, the index of the last stream element transferred is returned in Last. Otherwise, Item'First - 1 is returned in Last. Last is less than Item'Last only if the end of the stream is reached.
- 9 The Write operation appends Item to the specified stream.

Implementation Permissions

- 9.1/1 If Stream_Element'Size is not a multiple of System.Storage_Unit, then the components of Stream_Element_Array need not be aliased.

NOTES

31 See A.12.1, “The Package Streams.Stream_IO” for an example of extending type Root_Stream_Type.

11/2 32 If the end of stream has been reached, and Item'First is Stream_Element_Offset'First, Read will raise Constraint_Error.

13.13.2 Stream-Oriented Attributes

- 1/1 The operational attributes Write, Read, Output, and Input convert values to a stream of elements and reconstruct values from a stream.

Static Semantics

- 1.1/2 For every subtype S of an elementary type T, the following representation attribute is defined:

- 1.2/2 **S'Stream_Size** Denotes the number of bits occupied in a stream by items of subtype S. Hence, the number of stream elements required per item of elementary type T is:

$$T'Stream_Size / \text{Ada.Streams.Stream_Element}'Size$$
- 1.4/2 The value of this attribute is of type *universal_integer* and is a multiple of Stream_Element'Size.
- 1.5/2 Stream_Size may be specified for first subtypes via an **attribute_definition_clause**; the expression of such a clause shall be static, nonnegative, and a multiple of Stream_Element'Size.

Implementation Advice

- 1.6/2 If not specified, the value of Stream_Size for an elementary type should be the number of bits that corresponds to the minimum number of stream elements required by the first subtype of the type, rounded up to the nearest factor or multiple of the word size that is also a multiple of the stream element size.
- 1.7/2 The recommended level of support for the Stream_Size attribute is:
- 1.8/2 • A Stream_Size clause should be supported for a discrete or fixed point type T if the specified Stream_Size is a multiple of Stream_Element'Size and is no less than the size of the first subtype of T, and no greater than the size of the largest type of the same elementary class (signed integer, modular integer, enumeration, ordinary fixed point, or decimal fixed point).

Static Semantics

For every subtype S of a specific type T, the following attributes are defined.

S'Write S'Write denotes a procedure with the following specification:

```
procedure S'Write(
  Stream : not null access Ada.Streams.Root_Stream_Type'Class;
  Item   : in T)
```

S'Write writes the value of Item to Stream.

S'Read S'Read denotes a procedure with the following specification:

```
procedure S'Read(
  Stream : not null access Ada.Streams.Root_Stream_Type'Class;
  Item   : out T)
```

S'Read reads the value of Item from Stream.

For an untagged derived type, the Write (resp. Read) attribute is inherited according to the rules given in 13.1 if the attribute is available for the parent type at the point where T is declared. For a tagged derived type, these attributes are not inherited, but rather the default implementations are used.

The default implementations of the Write and Read attributes, where available, execute as follows:

For elementary types, Read reads (and Write writes) the number of stream elements implied by the Stream_Size for the type T; the representation of those stream elements is implementation defined. For composite types, the Write or Read attribute for each component is called in canonical order, which is last dimension varying fastest for an array, and positional aggregate order for a record. Bounds are not included in the stream if T is an array type. If T is a discriminated type, discriminants are included only if they have defaults. If T is a tagged type, the tag is not included. For type extensions, the Write or Read attribute for the parent type is called, followed by the Write or Read attribute of each component of the extension part, in canonical order. For a limited type extension, if the attribute of the parent type or any progenitor type of T is available anywhere within the immediate scope of T, and the attribute of the parent type or the type of any of the extension components is not available at the freezing point of T, then the attribute of T shall be directly specified.

Constraint_Error is raised by the predefined Write attribute if the value of the elementary item is outside the range of values representable using Stream_Size bits. For a signed integer type, an enumeration type, or a fixed point type, the range is unsigned only if the integer code for the lower bound of the first subtype is nonnegative, and a (symmetric) signed range that covers all values of the first subtype would require more than Stream_Size bits; otherwise the range is signed.

For every subtype S'Class of a class-wide type T'Class:

S'Class'Write

S'Class'Write denotes a procedure with the following specification:

```
procedure S'Class'Write(
  Stream : not null access Ada.Streams.Root_Stream_Type'Class;
  Item   : in T'Class)
```

Dispatches to the subprogram denoted by the Write attribute of the specific type identified by the tag of Item.

S'Class'Read S'Class'Read denotes a procedure with the following specification:

```
procedure S'Class'Read(
  Stream : not null access Ada.Streams.Root_Stream_Type'Class;
  Item   : out T'Class)
```

- 16 Dispatches to the subprogram denoted by the Read attribute of the specific type identified by the tag of Item.

Implementation Advice

- 17/2 This paragraph was deleted.

Static Semantics

- 18 For every subtype S of a specific type T, the following attributes are defined.

19 S'Output S'Output denotes a procedure with the following specification:

```
procedure S'Output(
    Stream : not null access Ada.Streams.Root_Stream_Type'Class;
    Item   : in T)
```

21 S'Output writes the value of Item to Stream, including any bounds or discriminants.

22 S'Input S'Input denotes a function with the following specification:

```
function S'Input(
    Stream : not null access Ada.Streams.Root_Stream_Type'Class)
    return T
```

24 S'Input reads and returns one value from Stream, using any bounds or discriminants written by a corresponding S'Output to determine how much to read.

25/2 For an untagged derived type, the Output (resp. Input) attribute is inherited according to the rules given in 13.1 if the attribute is available for the parent type at the point where T is declared. For a tagged derived type, these attributes are not inherited, but rather the default implementations are used.

25.1/2 The default implementations of the Output and Input attributes, where available, execute as follows:

- If T is an array type, S'Output first writes the bounds, and S'Input first reads the bounds. If T has discriminants without defaults, S'Output first writes the discriminants (using S'Write for each), and S'Input first reads the discriminants (using S'Read for each).

27/2 • S'Output then calls S'Write to write the value of Item to the stream. S'Input then creates an object (with the bounds or discriminants, if any, taken from the stream), passes it to S'Read, and returns the value of the object. Normal default initialization and finalization take place for this object (see 3.3.1, 7.6, and 7.6.1).

27.1/2 If T is an abstract type, then S'Input is an abstract function.

28 For every subtype S'Class of a class-wide type TClass:

29 S'Class'Output

S'Class'Output denotes a procedure with the following specification:

```
procedure S'Class'Output(
    Stream : not null access Ada.Streams.Root_Stream_Type'Class;
    Item   : in T'Class)
```

31/2 First writes the external tag of Item to Stream (by calling String'Output(Stream, Tags.External_Tag(Item'Tag)) — see 3.9) and then dispatches to the subprogram denoted by the Output attribute of the specific type identified by the tag. Tag_Error is raised if the tag of Item identifies a type declared at an accessibility level deeper than that of S.

32 S'Class'Input

S'Class'Input denotes a function with the following specification:

```
function S'Class'Input(
    Stream : not null access Ada.Streams.Root_Stream_Type'Class)
    return T'Class
```

First reads the external tag from *Stream* and determines the corresponding internal tag (by calling `Tags.Descendant_Tag(String'Input(Stream), S'Tag)` which might raise `Tag_Error` — see 3.9) and then dispatches to the subprogram denoted by the `Input` attribute of the specific type identified by the internal tag; returns that result. If the specific type identified by the internal tag is not covered by *TClass* or is abstract, `Constraint_Error` is raised.

34/2

In the default implementation of `Read` and `Input` for a composite type, for each scalar component that is a discriminant or whose `component_declaration` includes a `default_expression`, a check is made that the value returned by `Read` for the component belongs to its subtype. `Constraint_Error` is raised if this check fails. For other scalar components, no check is made. For each component that is of an access type, if the implementation can detect that the value returned by `Read` for the component is not a value of its subtype, `Constraint_Error` is raised. If the value is not a value of its subtype and this error is not detected, the component has an abnormal value, and erroneous execution can result (see 13.9.1). In the default implementation of `Read` for a composite type with defaulted discriminants, if the actual parameter of `Read` is constrained, a check is made that the discriminants read from the stream are equal to those of the actual parameter. `Constraint_Error` is raised if this check fails.

35/2

It is unspecified at which point and in which order these checks are performed. In particular, if `Constraint_Error` is raised due to the failure of one of these checks, it is unspecified how many stream elements have been read from the stream.

36/2

In the default implementation of `Read` and `Input` for a type, `End_Error` is raised if the end of the stream is reached before the reading of a value of the type is completed.

37/1

The stream-oriented attributes may be specified for any type via an `attribute_definition_clause`. The subprogram name given in such a clause shall not denote an abstract subprogram. Furthermore, if a stream-oriented attribute is specified for an interface type by an `attribute_definition_clause`, the subprogram name given in the clause shall statically denote a null procedure.

38/2

A stream-oriented attribute for a subtype of a specific type *T* is *available* at places where one of the following conditions is true:

39/2

- *T* is nonlimited.
- The `attribute_designator` is `Read` (resp. `Write`) and *T* is a limited record extension, and the attribute `Read` (resp. `Write`) is available for the parent type of *T* and for the types of all of the extension components.
- *T* is a limited untagged derived type, and the attribute was inherited for the type.
- The `attribute_designator` is `Input` (resp. `Output`), and *T* is a limited type, and the attribute `Read` (resp. `Write`) is available for *T*.
- The attribute has been specified via an `attribute_definition_clause`, and the `attribute_definition_clause` is visible.

40/2

41/2

42/2

43/2

44/2

A stream-oriented attribute for a subtype of a class-wide type *TClass* is available at places where one of the following conditions is true:

45/2

- *T* is nonlimited;
- the attribute has been specified via an `attribute_definition_clause`, and the `attribute_definition_clause` is visible; or
- the corresponding attribute of *T* is available, provided that if *T* has a partial view, the corresponding attribute is available at the end of the visible part where *T* is declared.

46/2

47/2

48/2

- 49/2 An **attribute_reference** for one of the stream-oriented attributes is illegal unless the attribute is available at the place of the **attribute_reference**. Furthermore, an **attribute_reference** for $T\text{Input}$ is illegal if T is an abstract type.
- 50/2 In the **parameter_and_result_profiles** for the stream-oriented attributes, the subtype of the Item parameter is the base subtype of T if T is a scalar type, and the first subtype otherwise. The same rule applies to the result of the Input attribute.
- 51/2 For an **attribute_definition_clause** specifying one of these attributes, the subtype of the Item parameter shall be the base subtype if scalar, and the first subtype otherwise. The same rule applies to the result of the Input function.
- 52/2 A type is said to *support external streaming* if Read and Write attributes are provided for sending values of such a type between active partitions, with Write marshalling the representation, and Read unmarshalling the representation. A limited type supports external streaming only if it has available Read and Write attributes. A type with a part that is of an access type supports external streaming only if that access type or the type of some part that includes the access type component, has Read and Write attributes that have been specified via an **attribute_definition_clause**, and that **attribute_definition_clause** is visible. An anonymous access type does not support external streaming. All other types support external streaming.

Erroneous Execution

- 53/2 If the internal tag returned by Descendant_Tag to $T'\text{Class}'\text{Input}$ identifies a type that is not library-level and whose tag has not been created, or does not exist in the partition at the time of the call, execution is erroneous.

Implementation Requirements

- 54/1 For every subtype S of a language-defined nonlimited specific type T , the output generated by $S'\text{Output}$ or $S'\text{Write}$ shall be readable by $S'\text{Input}$ or $S'\text{Read}$, respectively. This rule applies across partitions if the implementation conforms to the Distributed Systems Annex.
- 55/2 If Constraint_Error is raised during a call to Read because of failure of one the above checks, the implementation must ensure that the discriminants of the actual parameter of Read are not modified.

Implementation Permissions

- 56/2 The number of calls performed by the predefined implementation of the stream-oriented attributes on the Read and Write operations of the stream type is unspecified. An implementation may take advantage of this permission to perform internal buffering. However, all the calls on the Read and Write operations of the stream type needed to implement an explicit invocation of a stream-oriented attribute must take place before this invocation returns. An explicit invocation is one appearing explicitly in the program text, possibly through a generic instantiation (see 12.3).

NOTES

- 57 33 For a definite subtype S of a type T , only $T'\text{Write}$ and $T'\text{Read}$ are needed to pass an arbitrary value of the subtype through a stream. For an indefinite subtype S of a type T , $T'\text{Output}$ and $T'\text{Input}$ will normally be needed, since $T'\text{Write}$ and $T'\text{Read}$ do not pass bounds, discriminants, or tags.
- 58 34 User-specified attributes of $S'\text{Class}$ are not inherited by other class-wide types descended from S .

*Examples**Example of user-defined Write attribute:*

```
procedure My_Write(
  Stream : not null access Ada.Streams.Root_Stream_Type'Class;
  Item   : My_Integer'Base);
for My_Integer'Write use My_Write;
```

13.14 Freezing Rules

This clause defines a place in the program text where each declared entity becomes “frozen.” A use of an entity, such as a reference to it by name, or (for a type) an expression of the type, causes freezing of the entity in some contexts, as described below. The Legality Rules forbid certain kinds of uses of an entity in the region of text where it is frozen.

The *freezing* of an entity occurs at one or more places (*freezing points*) in the program text where the representation for the entity has to be fully determined. Each entity is frozen from its first freezing point to the end of the program text (given the ordering of compilation units defined in 10.1.4).

The end of a **declarative_part**, **protected_body**, or a declaration of a library package or generic library package, causes *freezing* of each entity declared within it, except for incomplete types. A noninstance body other than a renames-as-body causes freezing of each entity declared before it within the same **declarative_part**.

A construct that (explicitly or implicitly) references an entity can cause the *freezing* of the entity, as defined by subsequent paragraphs. At the place where a construct causes freezing, each **name**, **expression**, **implicit_dereference**, or **range** within the construct causes freezing:

- The occurrence of a **generic_instantiation** causes freezing; also, if a parameter of the instantiation is defaulted, the **default_expression** or **default_name** for that parameter causes freezing.
- The occurrence of an **object_declaration** that has no corresponding completion causes freezing.
- The declaration of a record extension causes freezing of the parent subtype.
- The declaration of a record extension, interface type, task unit, or protected unit causes freezing of any progenitor types specified in the declaration.

A static expression causes freezing where it occurs. An object name or nonstatic expression causes freezing where it occurs, unless the name or expression is part of a **default_expression**, a **default_name**, or a per-object expression of a component’s **constraint**, in which case, the freezing occurs later as part of another construct.

An implicit call freezes the same entities that would be frozen by an explicit call. This is true even if the implicit call is removed via implementation permissions.

If an expression is implicitly converted to a type or subtype *T*, then at the place where the expression causes freezing, *T* is frozen.

The following rules define which entities are frozen at the place where a construct causes freezing:

- At the place where an expression causes freezing, the type of the expression is frozen, unless the expression is an enumeration literal used as a **discrete_choice** of the **array_aggregate** of an **enumeration_representation_clause**.

- 11 • At the place where a name causes freezing, the entity denoted by the name is frozen, unless the name is a prefix of an expanded name; at the place where an object name causes freezing, the nominal subtype associated with the name is frozen.
- 11.1/1 • At the place where an `implicit_dereference` causes freezing, the nominal subtype associated with the `implicit_dereference` is frozen.
- 12 • At the place where a `range` causes freezing, the type of the `range` is frozen.
- 13 • At the place where an `allocator` causes freezing, the designated subtype of its type is frozen. If the type of the `allocator` is a derived type, then all ancestor types are also frozen.
- 14 • At the place where a callable entity is frozen, each subtype of its profile is frozen. If the callable entity is a member of an entry family, the index subtype of the family is frozen. At the place where a function call causes freezing, if a parameter of the call is defaulted, the `default_expression` for that parameter causes freezing.
- 15 • At the place where a subtype is frozen, its type is frozen. At the place where a type is frozen, any expressions or names within the full type definition cause freezing; the first subtype, and any component subtypes, index subtypes, and parent subtype of the type are frozen as well. For a specific tagged type, the corresponding class-wide type is frozen as well. For a class-wide type, the corresponding specific type is frozen as well.
- 15.1/2 • At the place where a specific tagged type is frozen, the primitive subprograms of the type are frozen.

Legality Rules

- 16 The explicit declaration of a primitive subprogram of a tagged type shall occur before the type is frozen (see 3.9.2).
- 17 A type shall be completely defined before it is frozen (see 3.11.1 and 7.3).
- 18 The completion of a deferred constant declaration shall occur before the constant is frozen (see 7.4).
- 19/1 An operational or representation item that directly specifies an aspect of an entity shall appear before the entity is frozen (see 13.1).

Dynamic Semantics

- 20/2 The tag (see 3.9) of a tagged type T is created at the point where T is frozen.

Annex A (normative) Predefined Language Environment

This Annex contains the specifications of library units that shall be provided by every implementation. There are three root library units: Ada, Interfaces, and System; other library units are children of these:

1

2/2

Standard — A.1 Ada — A.2 Assertions — 11.4.2 Asynchronous_Task_Control — D.11 Calendar — 9.6 Arithmetic — 9.6.1 Formatting — 9.6.1 Time_Zones — 9.6.1 Characters — A.3.1 Conversions — A.3.4 Handling — A.3.2 Latin_1 — A.3.3 Command_Line — A.15 Complex_Text_IO — G.1.3 Containers — A.18.1 Doubly_Linked_Lists — A.18.3 Generic_Array_Sort — A.18.16 Generic_Constrained_Array_Sort — A.18.16 Hashed_Maps — A.18.5 Hashed_Sets — A.18.8 Indefinite_Doubly_Linked_Lists — A.18.11 Indefinite_Hashed_Maps — A.18.12 Indefinite_Hashed_Sets — A.18.14 Indefinite_Ordered_Maps — A.18.13 Indefinite_Ordered_Sets — A.18.15 Indefinite_Vectors — A.18.10 Ordered_Maps — A.18.6 Ordered_Sets — A.18.9 Vectors — A.18.2 Decimal — F.2 Direct_IO — A.8.4 Directories — A.16 Information — A.16 Dispatching — D.2.1 EDF — D.2.6 Round_Robin — D.2.5 Dynamic_Priorities — D.5	Standard (...continued) Ada (...continued) Environment_Variables — A.17 Exceptions — 11.4.1 Execution_Time — D.14 Group_Budgets — D.14.2 Timers — D.14.1 Finalization — 7.6 Float_Text_IO — A.10.9 Float_Wide_Text_IO — A.11 Float_Wide_Wide_Text_IO — A.11 Integer_Text_IO — A.10.8 Integer_Wide_Text_IO — A.11 Integer_Wide_Wide_Text_IO — A.11 Interrupts — C.3.2 Names — C.3.2 IO_Exceptions — A.13 Numerics — A.5 Complex_Arrays — G.3.2 Complex_Elementary_Functions — G.1.2 Complex_Types — G.1.1 Discrete_Random — A.5.2 Elementary_Functions — A.5.1 Float_Random — A.5.2 Generic_Complex_Arrays — G.3.2 Generic_Complex_Elementary_Functions — G.1.2 Generic_Complex_Types — G.1.1 Generic_Elementary_Functions — A.5.1 Generic_Real_Arrays — G.3.1 Real_Arrays — G.3.1 Real_Time — D.8 Timing_Events — D.15 Sequential_IO — A.8.1 Storage_IO — A.9 Streams — 13.13.1 Stream_IO — A.12.1
--	--

Standard (...continued)

Ada (...continued)

- Strings — A.4.1
 - Bounded — A.4.4
 - Hash — A.4.9
- Fixed — A.4.3
 - Hash — A.4.9
- Hash — A.4.9
- Maps — A.4.2
 - Constants — A.4.6
- Unbounded — A.4.5
 - Hash — A.4.9
- Wide_Bounded — A.4.7
 - Wide_Hash — A.4.7
- Wide_Fixed — A.4.7
 - Wide_Hash — A.4.7
- Wide_Hash — A.4.7
- Wide_Maps — A.4.7
 - Wide_Constants — A.4.7
- Wide_Unbounded — A.4.7
 - Wide_Hash — A.4.7
- Wide_Wide_Bounded — A.4.8
 - Wide_Wide_Hash — A.4.8
- Wide_Wide_Fixed — A.4.8
 - Wide_Wide_Hash — A.4.8
- Wide_Wide_Hash — A.4.8
- Wide_Wide_Maps — A.4.8
 - Wide_Wide_Constants — A.4.8
- Wide_Wide_Unbounded — A.4.8
 - Wide_Wide_Hash — A.4.8

- Synchronous_Task_Control — D.10
- Tags — 3.9
- Generic_Dispatching_Constructor — 3.9
- Task_Attributes — C.7.2
- Task_Identification — C.7.1
- Task_Termination — C.7.3

Standard (...continued)

Ada (...continued)

- Text_IO — A.10.1
 - Bounded_IO — A.10.11
 - Complex_IO — G.1.3
 - Editing — F.3.3
 - Text_Streams — A.12.2
 - Unbounded_IO — A.10.12
- Unchecked_Conversion — 13.9
- Unchecked_Deallocation — 13.11.2
- Wide_Characters — A.3.1
- Wide_Text_IO — A.11
 - Complex_IO — G.1.4
 - Editing — F.3.4
 - Text_Streams — A.12.3
 - Wide_Bounded_IO — A.11
 - Wide_Unbounded_IO — A.11
- Wide_Wide_Characters — A.3.1
- Wide_Wide_Text_IO — A.11
 - Complex_IO — G.1.5
 - Editing — F.3.5
 - Text_Streams — A.12.4
 - Wide_Wide_Bounded_IO — A.11
 - Wide_Wide_Unbounded_IO — A.11

Interfaces — B.2

- C — B.3
 - Pointers — B.3.2
 - Strings — B.3.1
- COBOL — B.4
- Fortran — B.5

System — 13.7

- Address_To_Access_Conversions — 13.7.2
- Machine_Code — 13.8
- RPC — E.5
- Storage_Elements — 13.7.1
- Storage_Pools — 13.11

Implementation Requirements

- 3/2 The implementation shall ensure that each language-defined subprogram is reentrant in the sense that concurrent calls on the same subprogram perform as specified, so long as all parameters that could be passed by reference denote nonoverlapping objects.

Implementation Permissions

- 4 The implementation may restrict the replacement of language-defined compilation units. The implementation may restrict children of language-defined library units (other than Standard).

A.1 The Package Standard

This clause outlines the specification of the package Standard containing all predefined identifiers in the language. The corresponding package body is not specified by the language.

The operators that are predefined for the types declared in the package Standard are given in comments since they are implicitly declared. Italics are used for pseudo-names of anonymous types (such as *root_real*) and for undefined information (such as *implementation-defined*).

Static Semantics

The library package Standard has the following declaration:

```

package Standard is
  pragma Pure(Standard);
  type Boolean is (False, True);
  -- The predefined relational operators for this type are as follows:
  -- function "=" (Left, Right : Boolean'Base) return Boolean;
  -- function "/=" (Left, Right : Boolean'Base) return Boolean;
  -- function "<" (Left, Right : Boolean'Base) return Boolean;
  -- function "<=" (Left, Right : Boolean'Base) return Boolean;
  -- function ">" (Left, Right : Boolean'Base) return Boolean;
  -- function ">=" (Left, Right : Boolean'Base) return Boolean;
  -- The predefined logical operators and the predefined logical
  -- negation operator are as follows:
  -- function "and" (Left, Right : Boolean'Base) return Boolean'Base;
  -- function "or" (Left, Right : Boolean'Base) return Boolean'Base;
  -- function "xor" (Left, Right : Boolean'Base) return Boolean'Base;
  -- function "not" (Right : Boolean'Base) return Boolean'Base;
  -- The integer type root_integer and the
  -- corresponding universal type universal_integer are predefined.
  type Integer is range implementation-defined;
  subtype Natural is Integer range 0 .. Integer'Last;
  subtype Positive is Integer range 1 .. Integer'Last;
  -- The predefined operators for type Integer are as follows:
  -- function "=" (Left, Right : Integer'Base) return Boolean;
  -- function "/=" (Left, Right : Integer'Base) return Boolean;
  -- function "<" (Left, Right : Integer'Base) return Boolean;
  -- function "<=" (Left, Right : Integer'Base) return Boolean;
  -- function ">" (Left, Right : Integer'Base) return Boolean;
  -- function ">=" (Left, Right : Integer'Base) return Boolean;
  -- function "+" (Right : Integer'Base) return Integer'Base;
  -- function "-" (Right : Integer'Base) return Integer'Base;
  -- function "abs" (Right : Integer'Base) return Integer'Base;
  -- function "+" (Left, Right : Integer'Base) return Integer'Base;
  -- function "-" (Left, Right : Integer'Base) return Integer'Base;
  -- function "*" (Left, Right : Integer'Base) return Integer'Base;
  -- function "/" (Left, Right : Integer'Base) return Integer'Base;
  -- function "rem" (Left, Right : Integer'Base) return Integer'Base;
  -- function "mod" (Left, Right : Integer'Base) return Integer'Base;
  -- function "***" (Left : Integer'Base; Right : Natural)
  --               return Integer'Base;

```

```

19      -- The specification of each operator for the type
-- root_integer, or for any additional predefined integer
-- type, is obtained by replacing Integer by the name of the type
-- in the specification of the corresponding operator of the type
-- Integer. The right operand of the exponentiation operator
-- remains as subtype Natural.

20/2    -- The floating point type root_real and the
-- corresponding universal type universal_real are predefined.

21      type Float is digits implementation-defined;
22      -- The predefined operators for this type are as follows:
23
24      -- function "=" (Left, Right : Float) return Boolean;
-- function "/=" (Left, Right : Float) return Boolean;
-- function "<" (Left, Right : Float) return Boolean;
-- function "<=" (Left, Right : Float) return Boolean;
-- function ">" (Left, Right : Float) return Boolean;
-- function ">=" (Left, Right : Float) return Boolean;
25
26      -- function "+" (Right : Float) return Float;
-- function "-" (Right : Float) return Float;
-- function "abs" (Right : Float) return Float;
27
28      -- function "+" (Left, Right : Float) return Float;
-- function "-" (Left, Right : Float) return Float;
-- function "*" (Left, Right : Float) return Float;
-- function "/" (Left, Right : Float) return Float;
29
30      -- function "***" (Left : Float; Right : Integer'Base) return Float;

31      -- The specification of each operator for the type root_real, or for
-- any additional predefined floating point type, is obtained by
-- replacing Float by the name of the type in the specification of the
-- corresponding operator of the type Float.

32      -- In addition, the following operators are predefined for the root
-- numeric types:
33
34      function "*" (Left : root_integer; Right : root_real)
-- return root_real;
35      function "*" (Left : root_real; Right : root_integer)
-- return root_real;
36      function "/" (Left : root_real; Right : root_integer)
-- return root_real;
37
38      -- The type universal_fixed is predefined.
-- The only multiplying operators defined between
-- fixed point types are
39
40      function "***" (Left : universal_fixed; Right : universal_fixed)
-- return universal_fixed;
41      function "/" (Left : universal_fixed; Right : universal_fixed)
-- return universal_fixed;
42
43      -- The type universal_access is predefined.
-- The following equality operators are predefined:
44
45      function "=" (Left, Right : universal_access) return Boolean;
function "/=" (Left, Right : universal_access) return Boolean;

```

-- The declaration of type Character is based on the standard ISO 8859-1 character set.

35/2

-- There are no character literals corresponding to the positions for control characters.

-- They are indicated in italics in this definition. See 3.5.2.

-- The predefined operators for the type Character are the same as for

-- any enumeration type.

-- The declaration of type `Wide_Character` is based on the standard ISO/IEC 10646:2003 BMP character set. The first 256 positions have the same contents as type `Character`. See 3.5.2.

36

```
type Wide_Character is (nul, soh ... Hex_0000FFFE, Hex_0000FFFF);
```

36.2/2 -- The declaration of type `Wide_Wide_Character` is based on the full
 -- ISO/IEC 10646:2003 character set. The first 65536 positions have the
 -- same contents as type `Wide_Character`. See 3.5.2.

```
type Wide_Wide_Character is (nul, soh ... Hex_7FFFFFFE, Hex_7FFFFFFF);
for Wide_Wide_Character'Size use 32;
```

36.3/2 package ASCII is ... end ASCII; -- Obsolescent; see J.5

37 -- Predefined string types:

```
type String is array(Positive range <>) of Character;
pragma Pack(String);
```

38 -- The predefined operators for this type are as follows:

```
-- function "=" (Left, Right: String) return Boolean;
-- function "/=" (Left, Right: String) return Boolean;
-- function "<" (Left, Right: String) return Boolean;
-- function "<=" (Left, Right: String) return Boolean;
-- function ">" (Left, Right: String) return Boolean;
-- function ">=" (Left, Right: String) return Boolean;
```

40 -- function "&" (Left: String; Right: String) return String;
-- function "&" (Left: Character; Right: String) return String;
-- function "&" (Left: String; Right: Character) return String;
-- function "&" (Left: Character; Right: Character) return String;

41 type Wide_String is array(Positive range <>) of Wide_Character;
pragma Pack(Wide_String);

42 -- The predefined operators for this type correspond to those for String.

42.1/2 type Wide_Wide_String is array (Positive range <>)
 of Wide_Wide_Character;
pragma Pack (Wide_Wide_String);

42.2/2 -- The predefined operators for this type correspond to those for String.

43 type Duration is delta implementation-defined range implementation-defined;

44 -- The predefined operators for the type Duration are the same as for
 -- any fixed point type.

45 -- The predefined exceptions:

```
Constraint_Error: exception;
Program_Error : exception;
Storage_Error : exception;
Tasking_Error : exception;
```

47 end Standard;

48 Standard has no private part.

49/2 In each of the types `Character`, `Wide_Character`, and `Wide_Wide_Character`, the character literals for the space character (position 32) and the non-breaking space character (position 160) correspond to different values. Unless indicated otherwise, each occurrence of the character literal '' in this International Standard refers to the space character. Similarly, the character literals for hyphen (position 45) and soft hyphen (position 173) correspond to different values. Unless indicated otherwise, each occurrence of the character literal '-' in this International Standard refers to the hyphen character.

Dynamic Semantics

50 Elaboration of the body of Standard has no effect.

Implementation Permissions

51 An implementation may provide additional predefined integer types and additional predefined floating point types. Not all of these types need have names.

Implementation Advice

If an implementation provides additional named predefined integer types, then the names should end with “Integer” as in “Long_Integer”. If an implementation provides additional named predefined floating point types, then the names should end with “Float” as in “Long_Float”. 52

NOTES

1 Certain aspects of the predefined entities cannot be completely described in the language itself. For example, although the enumeration type Boolean can be written showing the two enumeration literals False and True, the short-circuit control forms cannot be expressed in the language. 53

2 As explained in 8.1, “Declarative Region” and 10.1.4, “The Compilation Process”, the declarative region of the package Standard encloses every library unit and consequently the main subprogram; the declaration of every library unit is assumed to occur within this declarative region. Library_items are assumed to be ordered in such a way that there are no forward semantic dependences. However, as explained in 8.3, “Visibility”, the only library units that are visible within a given compilation unit are the library units named by all with_clauses that apply to the given unit, and moreover, within the declarative region of a given library unit, that library unit itself. 54

3 If all block_statements of a program are named, then the name of each program unit can always be written as an expanded name starting with Standard (unless Standard is itself hidden). The name of a library unit cannot be a homograph of a name (such as Integer) that is already declared in Standard. 55

4 The exception Standard.Numeric_Error is defined in J.6. 56

A.2 The Package Ada*Static Semantics*

The following language-defined library package exists: 1

```
package Ada is
  pragma Pure(Ada);
end Ada;
```

Ada serves as the parent of most of the other language-defined library units; its declaration is empty (except for the pragma Pure). 3

Legality Rules

In the standard mode, it is illegal to compile a child of package Ada. 4

A.3 Character Handling

This clause presents the packages related to character processing: an empty pure package Characters and child packages Characters.Handling and Characters.Latin_1. The package Characters.Handling provides classification and conversion functions for Character data, and some simple functions for dealing with Wide_Character and Wide_Wide_Character data. The child package Characters.Latin_1 declares a set of constants initialized to values of type Character. 1/2

A.3.1 The Packages Characters, Wide_Characters, and Wide_Wide_Characters

Static Semantics

- 1 The library package Characters has the following declaration:

```
2 package Ada.Characters is
  pragma Pure(Characters);
end Ada.Characters;
```

- 3/2 The library package Wide_Characters has the following declaration:

```
4/2 package Ada.Wide_Characters is
  pragma Pure(Wide_Characters);
end Ada.Wide_Characters;
```

- 5/2 The library package Wide_Wide_Characters has the following declaration:

```
6/2 package Ada.Wide_Wide_Characters is
  pragma Pure(Wide_Wide_Characters);
end Ada.Wide_Wide_Characters;
```

Implementation Advice

- 7/2 If an implementation chooses to provide implementation-defined operations on Wide_Character or Wide_String (such as case mapping, classification, collating and sorting, etc.) it should do so by providing child units of Wide_Characters. Similarly if it chooses to provide implementation-defined operations on Wide_Wide_Character or Wide_Wide_String it should do so by providing child units of Wide_Wide_Characters.

A.3.2 The Package Characters.Handling

Static Semantics

- 1 The library package Characters.Handling has the following declaration:

```
2/2 with Ada.Characters.Conversions;
package Ada.Characters.Handling is
  pragma Pure(Handling);
  --Character classification functions
  function Is_Control          (Item : in Character) return Boolean;
  function Is_Graphic           (Item : in Character) return Boolean;
  function Is_Letter             (Item : in Character) return Boolean;
  function Is_Lower              (Item : in Character) return Boolean;
  function Is_Upper              (Item : in Character) return Boolean;
  function Is_Basic              (Item : in Character) return Boolean;
  function Is_Digit              (Item : in Character) return Boolean;
  function Is_Decimal_Digit     (Item : in Character) return Boolean
    renames Is_Digit;
  function Is_Hexadecimal_Digit (Item : in Character) return Boolean;
  function Is_Alphanumeric       (Item : in Character) return Boolean;
  function Is_Special            (Item : in Character) return Boolean;
  --Conversion functions for Character and String
  function To_Lower (Item : in Character) return Character;
  function To_Upper (Item : in Character) return Character;
  function To_Basic  (Item : in Character) return Character;
  function To_Lower (Item : in String)   return String;
  function To_Upper (Item : in String)   return String;
  function To_Basic  (Item : in String)   return String;
```

```

-- Classifications of and conversions between Character and ISO 646          8
subtype ISO_646 is                                         9
  Character range Character'Val(0) .. Character'Val(127);
function Is_ISO_646 (Item : in Character) return Boolean;           10
function Is_ISO_646 (Item : in String)      return Boolean;
function To_ISO_646 (Item      : in Character;                   11
                      Substitute : in ISO_646 := ' ')
  return ISO_646;
function To_ISO_646 (Item      : in String;                     12
                      Substitute : in ISO_646 := ' ')
  return String;
-- The functions Is_Character, Is_String, To_Character, To_String, To_Wide_Character, 13/2
-- and To_Wide_String are obsolescent; see J.14.
Paragraphs 14 through 18 were deleted.                                     14
end Ada.Characters.Handling;                                              15

```

In the description below for each function that returns a Boolean result, the effect is described in terms of the conditions under which the value True is returned. If these conditions are not met, then the function returns False.

Each of the following classification functions has a formal Character parameter, Item, and returns a Boolean result.

Is_Control	True if Item is a control character. A <i>control character</i> is a character whose position is in one of the ranges 0..31 or 127..159.	22
Is_Graphic	True if Item is a graphic character. A <i>graphic character</i> is a character whose position is in one of the ranges 32..126 or 160..255.	23
Is_Letter	True if Item is a letter. A <i>letter</i> is a character that is in one of the ranges 'A'..'Z' or 'a'..'z', or whose position is in one of the ranges 192..214, 216..246, or 248..255.	24
Is_Lower	True if Item is a lower-case letter. A <i>lower-case letter</i> is a character that is in the range 'a'..'z', or whose position is in one of the ranges 223..246 or 248..255.	25
Is_Upper	True if Item is an upper-case letter. An <i>upper-case letter</i> is a character that is in the range 'A'..'Z' or whose position is in one of the ranges 192..214 or 216..222.	26
Is_Basic	True if Item is a basic letter. A <i>basic letter</i> is a character that is in one of the ranges 'A'..'Z' and 'a'..'z', or that is one of the following: 'Æ', 'æ', 'Ð', 'ð', 'Þ', 'þ', or 'ß'.	27
Is_Digit	True if Item is a decimal digit. A <i>decimal digit</i> is a character in the range '0'..'9'.	28
Is.Decimal_Digit	A renaming of Is_Digit.	29
Is_Hexadecimal_Digit	True if Item is a hexadecimal digit. A <i>hexadecimal digit</i> is a character that is either a decimal digit or that is in one of the ranges 'A' .. 'F' or 'a' .. 'f'.	30
Is_Alphanumeric	True if Item is an alphanumeric character. An <i>alphanumeric character</i> is a character that is either a letter or a decimal digit.	31
Is_Special	True if Item is a special graphic character. A <i>special graphic character</i> is a graphic character that is not alphanumeric.	32

Each of the names To_Lower, To_Upper, and To_Basic refers to two functions: one that converts from Character to Character, and the other that converts from String to String. The result of each Character-to-Character function is described below, in terms of the conversion applied to Item, its formal Character

parameter. The result of each String-to-String conversion is obtained by applying to each element of the function's String parameter the corresponding Character-to-Character conversion; the result is the null String if the value of the formal parameter is the null String. The lower bound of the result String is 1.

- 34 To_Lower Returns the corresponding lower-case value for Item if Is_Upper(Item), and returns Item otherwise.
- 35 To_Upper Returns the corresponding upper-case value for Item if Is_Lower(Item) and Item has an upper-case form, and returns Item otherwise. The lower case letters 'B' and 'Y' do not have upper case forms.
- 36 To_Basic Returns the letter corresponding to Item but with no diacritical mark, if Item is a letter but not a basic letter; returns Item otherwise.
- 37 The following set of functions test for membership in the ISO 646 character range, or convert between ISO 646 and Character.
- 38 Is_ISO_646 The function whose formal parameter, Item, is of type Character returns True if Item is in the subtype ISO_646.
- 39 Is_ISO_646 The function whose formal parameter, Item, is of type String returns True if Is_ISO_646(Item(I)) is True for each I in Item'Range.
- 40 To_ISO_646 The function whose first formal parameter, Item, is of type Character returns Item if Is_ISO_646(Item), and returns the Substitute ISO_646 character otherwise.
- 41 To_ISO_646 The function whose first formal parameter, Item, is of type String returns the String whose Range is 1..Item'Length and each of whose elements is given by To_ISO_646 of the corresponding element in Item.

Paragraphs 42 through 48 were deleted.

Implementation Advice

- 49/2 This paragraph was deleted.

NOTES

50 5 A basic letter is a letter without a diacritical mark.

51 6 Except for the hexadecimal digits, basic letters, and ISO_646 characters, the categories identified in the classification functions form a strict hierarchy:

- 52 — Control characters
- 53 — Graphic characters
 - 54 — Alphanumeric characters
 - 55 — Letters
 - 56 — Upper-case letters
 - 57 — Lower-case letters
 - 58 — Decimal digits
 - 59 — Special graphic characters

A.3.3 The Package Characters.Latin_1

The package Characters.Latin_1 declares constants for characters in ISO 8859-1.

Static Semantics

The library package Characters.Latin_1 has the following declaration:

```

package Ada.Characters.Latin_1 is
  pragma Pure(Latin_1);

  -- Control characters:
    NUL          : constant Character := Character'Val(0);
    SOH          : constant Character := Character'Val(1);
    STX          : constant Character := Character'Val(2);
    ETX          : constant Character := Character'Val(3);
    EOT          : constant Character := Character'Val(4);
    ENQ          : constant Character := Character'Val(5);
    ACK          : constant Character := Character'Val(6);
    BEL          : constant Character := Character'Val(7);
    BS           : constant Character := Character'Val(8);
    HT           : constant Character := Character'Val(9);
    LF           : constant Character := Character'Val(10);
    VT           : constant Character := Character'Val(11);
    FF           : constant Character := Character'Val(12);
    CR           : constant Character := Character'Val(13);
    SO           : constant Character := Character'Val(14);
    SI           : constant Character := Character'Val(15);

    DLE          : constant Character := Character'Val(16);
    DC1          : constant Character := Character'Val(17);
    DC2          : constant Character := Character'Val(18);
    DC3          : constant Character := Character'Val(19);
    DC4          : constant Character := Character'Val(20);
    NAK          : constant Character := Character'Val(21);
    SYN          : constant Character := Character'Val(22);
    ETB          : constant Character := Character'Val(23);
    CAN          : constant Character := Character'Val(24);
    EM           : constant Character := Character'Val(25);
    SUB          : constant Character := Character'Val(26);
    ESC          : constant Character := Character'Val(27);
    FS           : constant Character := Character'Val(28);
    GS           : constant Character := Character'Val(29);
    RS           : constant Character := Character'Val(30);
    US           : constant Character := Character'Val(31);

  -- ISO 646 graphic characters:
    Space        : constant Character := ' ' ;   -- Character'Val(32)
    Exclamation : constant Character := '!' ;   -- Character'Val(33)
    Quotation    : constant Character := '"' ;   -- Character'Val(34)
    Number_Sign  : constant Character := '#' ;   -- Character'Val(35)
    Dollar_Sign  : constant Character := '$' ;   -- Character'Val(36)
    Percent_Sign : constant Character := '%' ;  -- Character'Val(37)
    Ampersand    : constant Character := '&' ;   -- Character'Val(38)
    Apostrophe   : constant Character := '\'' ;  -- Character'Val(39)
    Left_Parenthesis : constant Character := '(' ; -- Character'Val(40)
    Right_Parenthesis : constant Character := ')' ; -- Character'Val(41)
    Asterisk    : constant Character := '*' ;   -- Character'Val(42)
    Plus_Sign   : constant Character := '+' ;   -- Character'Val(43)
    Comma        : constant Character := ',' ;   -- Character'Val(44)
    Hyphen       : constant Character := '-' ;   -- Character'Val(45)
    Minus_Sign   : Character renames Hyphen;
    Full_Stop    : constant Character := '.' ;   -- Character'Val(46)
    Solidus     : constant Character := '/' ;   -- Character'Val(47)

```

```

9      -- Decimal digits '0' through '9' are at positions 48 through 57
10     Colon          : constant Character := ':';    -- Character'Val(58)
11     Semicolon      : constant Character := ';';    -- Character'Val(59)
12     Less_Than_Sign : constant Character := '<';   -- Character'Val(60)
13     Equals_Sign    : constant Character := '=';   -- Character'Val(61)
14     Greater_Than_Sign : constant Character := '>';  -- Character'Val(62)
15     Question       : constant Character := '?';   -- Character'Val(63)
16     Commercial_At : constant Character := '@';   -- Character'Val(64)

-- Letters 'A' through 'Z' are at positions 65 through 90
17     Left_Square_Bracket : constant Character := '['; -- Character'Val(91)
18     Reverse_Solidus     : constant Character := '\'; -- Character'Val(92)
19     Right_Square_Bracket : constant Character := ']'; -- Character'Val(93)
20     Circumflex         : constant Character := '^'; -- Character'Val(94)
21     Low_Line           : constant Character := '_'; -- Character'Val(95)

22     Grave            : constant Character := '`'; -- Character'Val(96)
23     LC_A              : constant Character := 'a'; -- Character'Val(97)
24     LC_B              : constant Character := 'b'; -- Character'Val(98)
25     LC_C              : constant Character := 'c'; -- Character'Val(99)
26     LC_D              : constant Character := 'd'; -- Character'Val(100)
27     LC_E              : constant Character := 'e'; -- Character'Val(101)
28     LC_F              : constant Character := 'f'; -- Character'Val(102)
29     LC_G              : constant Character := 'g'; -- Character'Val(103)
30     LC_H              : constant Character := 'h'; -- Character'Val(104)
31     LC_I              : constant Character := 'i'; -- Character'Val(105)
32     LC_J              : constant Character := 'j'; -- Character'Val(106)
33     LC_K              : constant Character := 'k'; -- Character'Val(107)
34     LC_L              : constant Character := 'l'; -- Character'Val(108)
35     LC_M              : constant Character := 'm'; -- Character'Val(109)
36     LC_N              : constant Character := 'n'; -- Character'Val(110)
37     LC_O              : constant Character := 'o'; -- Character'Val(111)

38     LC_P              : constant Character := 'p'; -- Character'Val(112)
39     LC_Q              : constant Character := 'q'; -- Character'Val(113)
40     LC_R              : constant Character := 'r'; -- Character'Val(114)
41     LC_S              : constant Character := 's'; -- Character'Val(115)
42     LC_T              : constant Character := 't'; -- Character'Val(116)
43     LC_U              : constant Character := 'u'; -- Character'Val(117)
44     LC_V              : constant Character := 'v'; -- Character'Val(118)
45     LC_W              : constant Character := 'w'; -- Character'Val(119)
46     LC_X              : constant Character := 'x'; -- Character'Val(120)
47     LC_Y              : constant Character := 'y'; -- Character'Val(121)
48     LC_Z              : constant Character := 'z'; -- Character'Val(122)
49     Left_Curly_Bracket : constant Character := '{'; -- Character'Val(123)
50     Vertical_Line      : constant Character := '|'; -- Character'Val(124)
51     Right_Curly_Bracket : constant Character := '}'; -- Character'Val(125)
52     Tilde              : constant Character := '~'; -- Character'Val(126)
53     DEL                : constant Character := Character'Val(127);

-- ISO 6429 control characters:
54     IS4                : Character renames FS;
55     IS3                : Character renames GS;
56     IS2                : Character renames RS;
57     IS1                : Character renames US;

```

```

Reserved_128          : constant Character := Character'Val(128);      17
Reserved_129          : constant Character := Character'Val(129);
BPH                  : constant Character := Character'Val(130);
NBH                  : constant Character := Character'Val(131);
Reserved_132          : constant Character := Character'Val(132);
NEL                  : constant Character := Character'Val(133);
SSA                  : constant Character := Character'Val(134);
ESA                  : constant Character := Character'Val(135);
HTS                  : constant Character := Character'Val(136);
HTJ                  : constant Character := Character'Val(137);
VTS                  : constant Character := Character'Val(138);
PLD                  : constant Character := Character'Val(139);
PLU                  : constant Character := Character'Val(140);
RI                   : constant Character := Character'Val(141);
SS2                  : constant Character := Character'Val(142);
SS3                  : constant Character := Character'Val(143);

DCS                  : constant Character := Character'Val(144);      18
PU1                  : constant Character := Character'Val(145);
PU2                  : constant Character := Character'Val(146);
STS                  : constant Character := Character'Val(147);
CCH                  : constant Character := Character'Val(148);
MW                   : constant Character := Character'Val(149);
SPA                  : constant Character := Character'Val(150);
EPA                  : constant Character := Character'Val(151);

SOS                  : constant Character := Character'Val(152);      19
Reserved_153          : constant Character := Character'Val(153);
SCI                  : constant Character := Character'Val(154);
CSI                  : constant Character := Character'Val(155);
ST                   : constant Character := Character'Val(156);
OSC                  : constant Character := Character'Val(157);
PM                   : constant Character := Character'Val(158);
APC                  : constant Character := Character'Val(159);

-- Other graphic characters:                                     20
-- Character positions 160 (16#A0#).. 175 (16#AF#):           21
No_Break_Space        : constant Character := ' ' ; --Character'Val(160)
NBSP                 : Character renames No_Break_Space;
Inverted_Exclamation : constant Character := '¡' ; --Character'Val(161)
Cent_Sign              : constant Character := '¢' ; --Character'Val(162)
Pound_Sign             : constant Character := '£' ; --Character'Val(163)
Currency_Sign          : constant Character := '¤' ; --Character'Val(164)
Yen_Sign               : constant Character := '¥' ; --Character'Val(165)
Broken_Bar             : constant Character := '׀' ; --Character'Val(166)
Section_Sign            : constant Character := '§' ; --Character'Val(167)
Diaeresis              : constant Character := '΅' ; --Character'Val(168)
Copyright_Sign          : constant Character := '©' ; --Character'Val(169)
Feminine_Ordinal_Indicator : constant Character := 'ª' ; --Character'Val(170)
Left_Angle_Quotation    : constant Character := '«' ; --Character'Val(171)
Not_Sign                : constant Character := '¬' ; --Character'Val(172)
Soft_Hyphen              : constant Character := '‐' ; --Character'Val(173)
Registered_Trade_Mark_Sign : constant Character := '®' ; --Character'Val(174)
Macron                 : constant Character := '‐' ; --Character'Val(175)

```

22

```
-- Character positions 176 (16#B0#) .. 191 (16#BF#):
Degree_Sign           : constant Character := '°'; --Character'Val(176)
Ring_Above             : Character renames Degree_Sign;
Plus_Minus_Sign       : constant Character := '±'; --Character'Val(177)
Superscript_Two        : constant Character := '²'; --Character'Val(178)
Superscript_Three      : constant Character := '³'; --Character'Val(179)
Acute                  : constant Character := '́'; --Character'Val(180)
Micro_Sign              : constant Character := 'µ'; --Character'Val(181)
Pilcrow_Sign            : constant Character := '¶'; --Character'Val(182)
Paragraph_Sign          : Character renames Pilcrow_Sign;
Middle_Dot              : constant Character := '·'; --Character'Val(183)
Cedilla                 : constant Character := '¸'; --Character'Val(184)
Superscript_One         : constant Character := '¹'; --Character'Val(185)
Masculine_Ordinal_Indicator: constant Character := 'º'; --Character'Val(186)
Right_Angle_Quotation    : constant Character := '»'; --Character'Val(187)
Fraction_One_Quarter     : constant Character := '¼'; --Character'Val(188)
Fraction_One_Half        : constant Character := '½'; --Character'Val(189)
Fraction_Three_Quiarters: constant Character := '¾'; --Character'Val(190)
Inverted_Question        : constant Character := '¿'; --Character'Val(191)
```

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```
-- Character positions 192 (16#C0#) .. 207 (16#CF#):
UC_A_Grave             : constant Character := 'À'; --Character'Val(192)
UC_A_Acute               : constant Character := 'Á'; --Character'Val(193)
UC_A_Circumflex          : constant Character := 'Â'; --Character'Val(194)
UC_A_Tilde                : constant Character := 'Ã'; --Character'Val(195)
UC_A_Diaeresis            : constant Character := 'Ä'; --Character'Val(196)
UC_A_Ring                  : constant Character := 'Å'; --Character'Val(197)
UC_AE_Diphthong           : constant Character := 'Æ'; --Character'Val(198)
UC_C_Cedilla              : constant Character := 'Ç'; --Character'Val(199)
UC_E_Grave                : constant Character := 'È'; --Character'Val(200)
UC_E_Acute                  : constant Character := 'É'; --Character'Val(201)
UC_E_Circumflex             : constant Character := 'Ê'; --Character'Val(202)
UC_E_Diaeresis              : constant Character := 'Ë'; --Character'Val(203)
UC_I_Grave                  : constant Character := 'Í'; --Character'Val(204)
UC_I_Acute                  : constant Character := 'Í'; --Character'Val(205)
UC_I_Circumflex              : constant Character := 'Î'; --Character'Val(206)
UC_I_Diaeresis              : constant Character := 'Ï'; --Character'Val(207)
```

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```
-- Character positions 208 (16#D0#) .. 223 (16#DF#):
UC_Icelandic_Eth          : constant Character := 'Ð'; --Character'Val(208)
UC_N_Tilde                  : constant Character := 'Ñ'; --Character'Val(209)
UC_O_Grave                  : constant Character := 'Ò'; --Character'Val(210)
UC_O_Acute                  : constant Character := 'Ó'; --Character'Val(211)
UC_O_Circumflex              : constant Character := 'Ô'; --Character'Val(212)
UC_O_Tilde                  : constant Character := 'Õ'; --Character'Val(213)
UC_O_Diaeresis              : constant Character := 'Ö'; --Character'Val(214)
Multiplication_Sign          : constant Character := '×'; --Character'Val(215)
UC_O_Oblique_Stroke          : constant Character := 'Ø'; --Character'Val(216)
UC_U_Grave                  : constant Character := 'Ù'; --Character'Val(217)
UC_U_Acute                  : constant Character := 'Û'; --Character'Val(218)
UC_U_Circumflex              : constant Character := 'Ü'; --Character'Val(219)
UC_U_Diaeresis              : constant Character := 'Ü'; --Character'Val(220)
UC_Y_Acute                  : constant Character := 'Ý'; --Character'Val(221)
UC_Icelandic_Thorn            : constant Character := 'Þ'; --Character'Val(222)
LC_German_Sharp_S            : constant Character := 'ß'; --Character'Val(223)
```

```

-- Character positions 224 (16#E0#) .. 239 (16#EF#): 25
LC_A_Grave           : constant Character := 'à'; --Character'Val(224)
LC_A_Acute            : constant Character := 'á'; --Character'Val(225)
LC_A_Circumflex       : constant Character := 'â'; --Character'Val(226)
LC_A_Tilde             : constant Character := 'ã'; --Character'Val(227)
LC_A_Diaeresis        : constant Character := 'ä'; --Character'Val(228)
LC_A_Ring              : constant Character := 'å'; --Character'Val(229)
LC_AE_Diphthong        : constant Character := 'æ'; --Character'Val(230)
LC_C_Cedilla           : constant Character := 'ç'; --Character'Val(231)
LC_E_Grave             : constant Character := 'è'; --Character'Val(232)
LC_E_Acute              : constant Character := 'é'; --Character'Val(233)
LC_E_Circumflex         : constant Character := 'ê'; --Character'Val(234)
LC_E_Diaeresis          : constant Character := 'ë'; --Character'Val(235)
LC_I_Grave              : constant Character := 'ì'; --Character'Val(236)
LC_I_Acute              : constant Character := 'í'; --Character'Val(237)
LC_I_Circumflex         : constant Character := 'î'; --Character'Val(238)
LC_I_Diaeresis          : constant Character := 'ï'; --Character'Val(239)

-- Character positions 240 (16#F0#) .. 255 (16#FF#): 26
LC_Icelandic_Eth       : constant Character := 'ð'; --Character'Val(240)
LC_N_Tilde              : constant Character := 'ñ'; --Character'Val(241)
LC_O_Grave              : constant Character := 'ð'; --Character'Val(242)
LC_O_Acute              : constant Character := 'ó'; --Character'Val(243)
LC_O_Circumflex          : constant Character := 'ô'; --Character'Val(244)
LC_O_Tilde              : constant Character := 'õ'; --Character'Val(245)
LC_O_Diaeresis          : constant Character := 'ö'; --Character'Val(246)
Division_Sign           : constant Character := '+'; --Character'Val(247)
LC_O_Oblique_Stroke     : constant Character := 'ø'; --Character'Val(248)
LC_U_Grave              : constant Character := 'ù'; --Character'Val(249)
LC_U_Acute              : constant Character := 'ú'; --Character'Val(250)
LC_U_Circumflex          : constant Character := 'û'; --Character'Val(251)
LC_U_Diaeresis          : constant Character := 'ü'; --Character'Val(252)
LC_Y_Acute              : constant Character := 'ÿ'; --Character'Val(253)
LC_Icelandic_Thorn      : constant Character := 'þ'; --Character'Val(254)
LC_Y_Diaeresis          : constant Character := 'ÿ'; --Character'Val(255)

end Ada.Characters.Latin_1;

```

Implementation Permissions

An implementation may provide additional packages as children of Ada.Characters, to declare names for the symbols of the local character set or other character sets.

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A.3.4 The Package Characters.Conversions

Static Semantics

1/2 The library package Characters.Conversions has the following declaration:

```
2/2  package Ada.Characters.Conversions is
      pragma Pure(Conversions);

3/2    function Is_Character (Item : in Wide_Character)      return Boolean;
    function Is_String     (Item : in Wide_String)        return Boolean;
    function Is_Character (Item : in Wide_Wide_Character) return Boolean;
    function Is_String     (Item : in Wide_Wide_String)   return Boolean;
    function Is_Wide_Character (Item : in Wide_Wide_Character)
      return Boolean;
    function Is_Wide_String   (Item : in Wide_Wide_String)
      return Boolean;

4/2    function To_Wide_Character (Item : in Character)  return Wide_Character;
    function To_Wide_String   (Item : in String)       return Wide_String;
    function To_Wide_Wide_Character (Item : in Character)
      return Wide_Wide_Character;
    function To_Wide_Wide_String  (Item : in String)
      return Wide_Wide_String;
    function To_Wide_Wide_Character (Item : in Wide_Character)
      return Wide_Wide_Character;
    function To_Wide_Wide_String  (Item : in Wide_String)
      return Wide_Wide_String;

5/2    function To_Character (Item      : in Wide_Character;
                           Substitute : in Character := ' ')
      return Character;
    function To_String   (Item      : in Wide_String;
                           Substitute : in Character := ' ')
      return String;
    function To_Character (Item      : in Wide_Wide_Character;
                           Substitute : in Character := ' ')
      return Character;
    function To_String   (Item      : in Wide_Wide_String;
                           Substitute : in Character := ' ')
      return String;
    function To_Wide_Character (Item      : in Wide_Wide_Character;
                               Substitute : in Wide_Character := ' ')
      return Wide_Character;
    function To_Wide_String   (Item      : in Wide_Wide_String;
                               Substitute : in Wide_Character := ' ')
      return Wide_String;

6/2  end Ada.Characters.Conversions;
```

7/2 The functions in package Characters.Conversions test Wide_Wide_Character or Wide_Character values for membership in Wide_Character or Character, or convert between corresponding characters of Wide_Wide_Character, Wide_Character, and Character.

```
8/2    function Is_Character (Item : in Wide_Character) return Boolean;
9/2      Returns True if Wide_Character'Pos(Item) <= Character'Pos(Character'Last).

10/2   function Is_Character (Item : in Wide_Wide_Character) return Boolean;
11/2     Returns True if Wide_Wide_Character'Pos(Item) <= Character'Pos(Character'Last).

12/2   function Is_Wide_Character (Item : in Wide_Wide_Character) return Boolean;
13/2     Returns True if Wide_Wide_Character'Pos(Item) <= Wide_Character'Pos(Wide_Character'Last).
```

function Is_String (Item : in Wide_String) return Boolean;	14/2
function Is_String (Item : in Wide_Wide_String) return Boolean;	
>Returns True if Is_Character(Item(I)) is True for each I in Item'Range.	15/2
function Is_Wide_String (Item : in Wide_Wide_String) return Boolean;	16/2
>Returns True if Is_Wide_Character(Item(I)) is True for each I in Item'Range.	17/2
function To_Character (Item : in Wide_Character; Substitute : in Character := ' ') return Character;	18/2
function To_Character (Item : in Wide_Wide_Character; Substitute : in Character := ' ') return Character;	
>Returns the Character corresponding to Item if Is_Character(Item), and returns the Substitute Character otherwise.	19/2
function To_Wide_Character (Item : in Character) return Wide_Character;	20/2
>Returns the Wide_Character X such that Character'Pos(Item) = Wide_Character'Pos (X).	21/2
function To_Wide_Character (Item : in Wide_Wide_Character; Substitute : in Wide_Character := ' ') return Wide_Character;	22/2
>Returns the Wide_Character corresponding to Item if Is_Wide_Character(Item), and returns the Substitute Wide_Character otherwise.	23/2
function To_Wide_Wide_Character (Item : in Character) return Wide_Wide_Character;	24/2
>Returns the Wide_Wide_Character X such that Character'Pos(Item) = Wide_Wide_Character'Pos (X).	25/2
function To_Wide_Wide_Character (Item : in Wide_Character) return Wide_Wide_Character;	26/2
>Returns the Wide_Wide_Character X such that Wide_Character'Pos(Item) = Wide_Wide_Character'Pos (X).	27/2
function To_String (Item : in Wide_String; Substitute : in Character := ' ') return String;	28/2
function To_String (Item : in Wide_Wide_String; Substitute : in Character := ' ') return String;	
>Returns the String whose range is 1..Item'Length and each of whose elements is given by To_Character of the corresponding element in Item.	29/2
function To_Wide_String (Item : in String) return Wide_String;	30/2
>Returns the Wide_String whose range is 1..Item'Length and each of whose elements is given by To_Wide_Character of the corresponding element in Item.	31/2
function To_Wide_String (Item : in Wide_Wide_String; Substitute : in Wide_Character := ' ') return Wide_String;	32/2
>Returns the Wide_String whose range is 1..Item'Length and each of whose elements is given by To_Wide_Character of the corresponding element in Item with the given Substitute Wide_Character.	33/2

34/2 **function** To_Wide_Wide_String (Item : in String) **return** Wide_Wide_String;
 function To_Wide_Wide_String (Item : in Wide_String)
 return Wide_Wide_String;

35/2 Returns the Wide_Wide_String whose range is 1..Item'Length and each of whose elements is given by To_Wide_Wide_Character of the corresponding element in Item.

A.4 String Handling

1/2 This clause presents the specifications of the package Strings and several child packages, which provide facilities for dealing with string data. Fixed-length, bounded-length, and unbounded-length strings are supported, for String, Wide_String, and Wide_Wide_String. The string-handling subprograms include searches for pattern strings and for characters in program-specified sets, translation (via a character-to-character mapping), and transformation (replacing, inserting, overwriting, and deleting of substrings).

A.4.1 The Package Strings

1 The package Strings provides declarations common to the string handling packages.

Static Semantics

2 The library package Strings has the following declaration:

```
3       package Ada.Strings is
           pragma Pure(Strings);
4/2       Space       : constant Character       := ' ';
       Wide_Space   : constant Wide_Character   := ' ';
       Wide_Wide_Space   : constant Wide_Wide_Character   := ' ';
5       Length_Error, Pattern_Error, Index_Error, Translation_Error : exception;
6       type Alignment is (Left, Right, Center);
       type Truncation is (Left, Right, Error);
       type Membership is (Inside, Outside);
       type Direction is (Forward, Backward);
       type Trim_End is (Left, Right, Both);
end Ada.Strings;
```

A.4.2 The Package Strings.Maps

1 The package Strings.Maps defines the types, operations, and other entities needed for character sets and character-to-character mappings.

Static Semantics

2 The library package Strings.Maps has the following declaration:

```
3/2       package Ada.Strings.Maps is
           pragma Pure(Maps);
4/2       -- Representation for a set of character values:
       type Character_Set is private;
       pragma Preelaborable_Initialization(Character_Set);
5       Null_Set : constant Character_Set;
6       type Character_Range is
           record
              Low   : Character;
              High   : Character;
           end record;
       -- Represents Character range Low..High
```

```

type Character_Ranges is array (Positive range <>) of Character_Range;          7
function To_Set      (Ranges : in Character_Ranges) return Character_Set;        8
function To_Set      (Span   : in Character_Range)  return Character_Set;        9
function To_Ranges   (Set    : in Character_Set)   return Character_Ranges;       10
function "="        (Left, Right : in Character_Set) return Boolean;            11
function "not"      (Right : in Character_Set)   return Character_Set;          12
function "and"      (Left, Right : in Character_Set) return Character_Set;
function "or"       (Left, Right : in Character_Set) return Character_Set;
function "xor"      (Left, Right : in Character_Set) return Character_Set;
function "-"        (Left, Right : in Character_Set) return Character_Set;
function Is_In      (Element : in Character;
                     Set     : in Character_Set)
                     return Boolean;                           13
function Is_Subset  (Elements : in Character_Set;
                     Set     : in Character_Set)
                     return Boolean;                           14
function "<="       (Left   : in Character_Set;
                     Right  : in Character_Set)
                     return Boolean renames Is_Subset;        15
-- Alternative representation for a set of character values:
subtype Character_Sequence is String;
function To_Set     (Sequence : in Character_Sequence) return Character_Set;    17
function To_Set     (Singleton : in Character)      return Character_Set;        18
function To_Sequence (Set   : in Character_Set)  return Character_Sequence;     19
-- Representation for a character to character mapping:
type Character_Mapping is private;                                         20/2
pragma Preelaborable_Initialization(Character_Mapping);
function Value      (Map    : in Character_Mapping;
                     Element : in Character)
                     return Character;                         21
Identity : constant Character_Mapping;                                     22
function To_Mapping  (From, To : in Character_Sequence) return Character_Mapping; 23
function To_Domain   (Map : in Character_Mapping)
                     return Character_Sequence;                24
function To_Range    (Map : in Character_Mapping)
                     return Character_Sequence;               25
type Character_Mapping_Function is
  access function (From : in Character) return Character;
private
  ... -- not specified by the language
end Ada.Strings.Maps;                                              26

An object of type Character_Set represents a set of characters.                      27
Null_Set represents the set containing no characters.                            28
An object Obj of type Character_Range represents the set of characters in the range Obj.Low .. Obj.High. 29
An object Obj of type Character_Ranges represents the union of the sets corresponding to Obj(I) for I in Obj'Range. 30
function To_Set      (Ranges : in Character_Ranges) return Character_Set;        31
If Ranges'Length=0 then Null_Set is returned; otherwise the returned value represents the set      32
corresponding to Ranges.

```

```

33      function To_Set (Span : in Character_Range) return Character_Set;
34          The returned value represents the set containing each character in Span.
35      function To_Ranges (Set : in Character_Set) return Character_Ranges;
36          If Set = Null_Set then an empty Character_Ranges array is returned; otherwise the shortest array
            of contiguous ranges of Character values in Set, in increasing order of Low, is returned.
37      function "=" (Left, Right : in Character_Set) return Boolean;
38          The function "=" returns True if Left and Right represent identical sets, and False otherwise.
39
        Each of the logical operators "not", "and", "or", and "xor" returns a Character_Set value that represents
        the set obtained by applying the corresponding operation to the set(s) represented by the parameter(s) of
        the operator. "-"(Left, Right) is equivalent to "and"(Left, "not"(Right)).
40      function Is_In (Element : in Character;
41                      Set      : in Character_Set);
42          return Boolean;
43
        Is_In returns True if Element is in Set, and False otherwise.
44      function Is_Subset (Elements : in Character_Set;
45                          Set      : in Character_Set)
46          return Boolean;
47
        Is_Subset returns True if Elements is a subset of Set, and False otherwise.
48      subtype Character_Sequence is String;
49
        The Character_Sequence subtype is used to portray a set of character values and also to identify
        the domain and range of a character mapping.
50      function To_Set (Sequence : in Character_Sequence) return Character_Set;
51      function To_Set (Singleton : in Character)           return Character_Set;
52
        Sequence portrays the set of character values that it explicitly contains (ignoring duplicates).
        Singleton portrays the set comprising a single Character. Each of the To_Set functions returns a
        Character_Set value that represents the set portrayed by Sequence or Singleton.
53      function To_Sequence (Set : in Character_Set) return Character_Sequence;
54
        The function To_Sequence returns a Character_Sequence value containing each of the characters
        in the set represented by Set, in ascending order with no duplicates.
55      type Character_Mapping is private;
56
        An object of type Character_Mapping represents a Character-to-Character mapping.
57      function Value (Map      : in Character_Mapping;
58                      Element : in Character)
59          return Character;
60
        The function Value returns the Character value to which Element maps with respect to the
        mapping represented by Map.
61
        A character C matches a pattern character P with respect to a given Character_Mapping value Map if
        Value(Map, C) = P. A string S matches a pattern string P with respect to a given Character_Mapping if
        their lengths are the same and if each character in S matches its corresponding character in the pattern
        string P.

```

String handling subprograms that deal with character mappings have parameters whose type is Character_Mapping. 55

`Identity : constant Character_Mapping;` 56

Identity maps each Character to itself. 57

`function To_Mapping (From, To : in Character_Sequence)
return Character_Mapping;` 58

To_Mapping produces a Character_Mapping such that each element of From maps to the corresponding element of To, and each other character maps to itself. If From'Length /= To'Length, or if some character is repeated in From, then Translation_Error is propagated. 59

`function To_Domain (Map : in Character_Mapping) return Character_Sequence;` 60

To_Domain returns the shortest Character_Sequence value D such that each character not in D maps to itself, and such that the characters in D are in ascending order. The lower bound of D is 1. 61

`function To_Range (Map : in Character_Mapping) return Character_Sequence;` 62

To_Range returns the Character_Sequence value R, such that if D = To_Domain(Map), then R has the same bounds as D, and D(I) maps to R(I) for each I in D'Range. 63/1

An object F of type Character_Mapping_Function maps a Character value C to the Character value F.all(C), which is said to *match* C with respect to mapping function F. 64

NOTES

7 Character_Mapping and Character_Mapping_Function are used both for character equivalence mappings in the search subprograms (such as for case insensitivity) and as transformational mappings in the Translate subprograms. 65

8 To_Domain(Identity) and To_Range(Identity) each returns the null string. 66

Examples

To_Mapping("ABCD", "ZZAB") returns a Character_Mapping that maps 'A' and 'B' to 'Z', 'C' to 'A', 'D' to 'B', and each other Character to itself. 67

A.4.3 Fixed-Length String Handling

The language-defined package Strings.Fixed provides string-handling subprograms for fixed-length strings; that is, for values of type Standard.String. Several of these subprograms are procedures that modify the contents of a String that is passed as an **out** or an **in out** parameter; each has additional parameters to control the effect when the logical length of the result differs from the parameter's length. 1

For each function that returns a String, the lower bound of the returned value is 1. 2

The basic model embodied in the package is that a fixed-length string comprises significant characters and possibly padding (with space characters) on either or both ends. When a shorter string is copied to a longer string, padding is inserted, and when a longer string is copied to a shorter one, padding is stripped. The Move procedure in Strings.Fixed, which takes a String as an **out** parameter, allows the programmer to control these effects. Similar control is provided by the string transformation procedures. 3

Static Semantics

4 The library package Strings.Fixed has the following declaration:

```

5   with Ada.Strings.Maps;
6   package Ada.Strings.Fixed is
7     pragma Preelaborate(Fixed);
8   -- "Copy" procedure for strings of possibly different lengths
9     procedure Move (Source  : in String;
10                  Target   : out String;
11                  Drop    : in Truncation := Error;
12                  Justify : in Alignment  := Left;
13                  Pad     : in Character   := Space);
14
15   -- Search subprograms
16
17   function Index (Source  : in String;
18                  Pattern : in String;
19                  From    : in Positive;
20                  Going   : in Direction := Forward;
21                  Mapping  : in Maps.Character_Mapping := Maps.Identity)
22     return Natural;
23
24   function Index (Source  : in String;
25                  Pattern : in String;
26                  From    : in Positive;
27                  Going   : in Direction := Forward;
28                  Mapping  : in Maps.Character_Mapping_Function)
29     return Natural;
30
31   function Index (Source  : in String;
32                  Pattern : in String;
33                  Going   : in Direction := Forward;
34                  Mapping  : in Maps.Character_Mapping
35                               := Maps.Identity)
36     return Natural;
37
38   function Index (Source  : in String;
39                  Pattern : in String;
40                  Going   : in Direction := Forward;
41                  Mapping  : in Maps.Character_Mapping_Function)
42     return Natural;
43
44   function Index (Source  : in String;
45                  Set      : in Maps.Character_Set;
46                  From    : in Positive;
47                  Test    : in Membership := Inside;
48                  Going   : in Direction := Forward)
49     return Natural;
50
51   function Index (Source : in String;
52                  Set   : in Maps.Character_Set;
53                  Test  : in Membership := Inside;
54                  Going : in Direction := Forward)
55     return Natural;
56
57   function Index_Non_Bank (Source : in String;
58                           From   : in Positive;
59                           Going  : in Direction := Forward)
60     return Natural;
61
62   function Index_Non_Bank (Source : in String;
63                           Going  : in Direction := Forward)
64     return Natural;
65
66   function Count (Source  : in String;
67                  Pattern : in String;
68                  Mapping  : in Maps.Character_Mapping
69                               := Maps.Identity)
70     return Natural;

```

```

function Count (Source : in String;
                Pattern : in String;
                Mapping : in Maps.Character_Mapping_Function) 14
    return Natural;
function Count (Source : in String;
                Set : in Maps.Character_Set) 15
    return Natural;
procedure Find_Token (Source : in String;
                      Set : in Maps.Character_Set;
                      Test : in Membership;
                      First : out Positive;
                      Last : out Natural); 16
-- String translation subprograms 17
function Translate (Source : in String;
                     Mapping : in Maps.Character_Mapping) 18
    return String;
procedure Translate (Source : in out String;
                     Mapping : in Maps.Character_Mapping); 19
function Translate (Source : in String;
                     Mapping : in Maps.Character_Mapping_Function) 20
    return String;
procedure Translate (Source : in out String;
                     Mapping : in Maps.Character_Mapping_Function); 21
-- String transformation subprograms 22
function Replace_Slice (Source : in String;
                         Low : in Positive;
                         High : in Natural;
                         By : in String) 23
    return String;
procedure Replace_Slice (Source : in out String;
                         Low : in Positive;
                         High : in Natural;
                         By : in String;
                         Drop : in Truncation := Error;
                         Justify : in Alignment := Left;
                         Pad : in Character := Space); 24
function Insert (Source : in String;
                  Before : in Positive;
                  New_Item : in String) 25
    return String;
procedure Insert (Source : in out String;
                  Before : in Positive;
                  New_Item : in String;
                  Drop : in Truncation := Error); 26
function Overwrite (Source : in String;
                      Position : in Positive;
                      New_Item : in String) 27
    return String;
procedure Overwrite (Source : in out String;
                      Position : in Positive;
                      New_Item : in String;
                      Drop : in Truncation := Right); 28
function Delete (Source : in String;
                  From : in Positive;
                  Through : in Natural) 29
    return String;

```

```

30      procedure Delete (Source  : in out String;
                      From    : in Positive;
                      Through : in Natural;
                      Justify : in Alignment := Left;
                      Pad     : in Character := Space);
31      --String selector subprograms
32      function Trim (Source : in String;
                           Side   : in Trim_End)
                    return String;
33      procedure Trim (Source  : in out String;
                           Side   : in Trim_End;
                           Justify : in Alignment := Left;
                           Pad     : in Character := Space);
34      function Trim (Source : in String;
                           Left   : in Maps.Character_Set;
                           Right  : in Maps.Character_Set)
                    return String;
35      procedure Trim (Source  : in out String;
                           Left   : in Maps.Character_Set;
                           Right  : in Maps.Character_Set;
                           Justify : in Alignment := Strings.Left;
                           Pad     : in Character := Space);
36      function Head (Source : in String;
                           Count  : in Natural;
                           Pad    : in Character := Space)
                    return String;
37      procedure Head (Source  : in out String;
                           Count  : in Natural;
                           Justify : in Alignment := Left;
                           Pad    : in Character := Space);
38      function Tail (Source : in String;
                           Count  : in Natural;
                           Pad    : in Character := Space)
                    return String;
39      procedure Tail (Source  : in out String;
                           Count  : in Natural;
                           Justify : in Alignment := Left;
                           Pad    : in Character := Space);
40      --String constructor functions
41      function "*" (Left  : in Natural;
                      Right : in Character) return String;
42      function "*" (Left  : in Natural;
                      Right : in String) return String;
43  end Ada.Strings.Fixed;

```

The effects of the above subprograms are as follows.

```

44      procedure Move (Source  : in String;
                           Target  : out String;
                           Drop    : in Truncation := Error;
                           Justify : in Alignment := Left;
                           Pad     : in Character := Space);

```

The Move procedure copies characters from Source to Target. If Source has the same length as Target, then the effect is to assign Source to Target. If Source is shorter than Target then:

- If Justify=Left, then Source is copied into the first Source'Length characters of Target.
- If Justify=Right, then Source is copied into the last Source'Length characters of Target.

- If Justify=Center, then Source is copied into the middle Source'Length characters of Target. In this case, if the difference in length between Target and Source is odd, then the extra Pad character is on the right. 48
 - Pad is copied to each Target character not otherwise assigned. 49
- If Source is longer than Target, then the effect is based on Drop. 50
- If Drop=Left, then the rightmost Target'Length characters of Source are copied into Target. 51
 - If Drop=Right, then the leftmost Target'Length characters of Source are copied into Target. 52
 - If Drop=Error, then the effect depends on the value of the Justify parameter and also on whether any characters in Source other than Pad would fail to be copied:
 - If Justify=Left, and if each of the rightmost Source'Length-Target'Length characters in Source is Pad, then the leftmost Target'Length characters of Source are copied to Target. 54
 - If Justify=Right, and if each of the leftmost Source'Length-Target'Length characters in Source is Pad, then the rightmost Target'Length characters of Source are copied to Target. 55
 - Otherwise, Length_Error is propagated. 56

```
function Index (Source   : in String;
                Pattern  : in String;
                From     : in Positive;
                Going    : in Direction := Forward;
                Mapping   : in Maps.Character_Mapping := Maps.Identity)
return Natural;

function Index (Source   : in String;
                Pattern  : in String;
                From     : in Positive;
                Going    : in Direction := Forward;
                Mapping   : in Maps.Character_Mapping_Function)
return Natural;
```

Each Index function searches, starting from From, for a slice of Source, with length Pattern'Length, that matches Pattern with respect to Mapping; the parameter Going indicates the direction of the lookup. If From is not in Source'Range, then Index_Error is propagated. If Going = Forward, then Index returns the smallest index I which is greater than or equal to From such that the slice of Source starting at I matches Pattern. If Going = Backward, then Index returns the largest index I such that the slice of Source starting at I matches Pattern and has an upper bound less than or equal to From. If there is no such slice, then 0 is returned. If Pattern is the null string, then Pattern_Error is propagated. 56.2/2

```

57      function Index (Source    : in String;
                  Pattern   : in String;
                  Going     : in Direction := Forward;
                  Mapping   : in Maps.Character_Mapping
                                := Maps.Identity)
        return Natural;

      function Index (Source    : in String;
                  Pattern   : in String;
                  Going     : in Direction := Forward;
                  Mapping   : in Maps.Character_Mapping_Function)
        return Natural;

58/2    If Going = Forward, returns
58.1/2   Index (Source, Pattern, Source'First, Forward, Mapping);
58.2/2   otherwise returns
58.3/2   Index (Source, Pattern, Source'Last, Backward, Mapping);

58.4/2   function Index (Source  : in String;
                      Set      : in Maps.Character_Set;
                      From     : in Positive;
                      Test     : in Membership := Inside;
                      Going   : in Direction := Forward)
        return Natural;

58.5/2   Index searches for the first or last occurrence of any of a set of characters (when Test=Inside), or
any of the complement of a set of characters (when Test=Outside). If From is not in
Source'Range, then Index_Error is propagated. Otherwise, it returns the smallest index I >=
From (if Going=Forward) or the largest index I <= From (if Going=Backward) such that
Source(I) satisfies the Test condition with respect to Set; it returns 0 if there is no such Character
in Source.

59      function Index (Source : in String;
                      Set    : in Maps.Character_Set;
                      Test   : in Membership := Inside;
                      Going  : in Direction := Forward)
        return Natural;

60/2    If Going = Forward, returns
60.1/2   Index (Source, Set, Source'First, Test, Forward);
60.2/2   otherwise returns
60.3/2   Index (Source, Set, Source'Last, Test, Backward);

60.4/2   function Index_Non_Bank (Source : in String;
                                 From   : in Positive;
                                 Going  : in Direction := Forward)
        return Natural;

60.5/2   Returns Index (Source, Maps.To_Set(Space), From, Outside, Going);

61      function Index_Non_Bank (Source : in String;
                                 Going  : in Direction := Forward)
        return Natural;

62      Returns Index (Source, Maps.To_Set(Space), Outside, Going)

```

function Count (Source : in String; Pattern : in String; Mapping : in Maps.Character_Mapping := Maps.Identity)	63
return Natural;	
function Count (Source : in String; Pattern : in String; Mapping : in Maps.Character_Mapping_Function)	64
return Natural;	
Returns the maximum number of nonoverlapping slices of Source that match Pattern with respect to Mapping. If Pattern is the null string then Pattern_Error is propagated.	64
function Count (Source : in String; Set : in Maps.Character_Set)	65
return Natural;	
Returns the number of occurrences in Source of characters that are in Set.	66
procedure Find_Token (Source : in String; Set : in Maps.Character_Set; Test : in Membership; First : out Positive; Last : out Natural);	67
Find_Token returns in First and Last the indices of the beginning and end of the first slice of Source all of whose elements satisfy the Test condition, and such that the elements (if any) immediately before and after the slice do not satisfy the Test condition. If no such slice exists, then the value returned for Last is zero, and the value returned for First is Source'First; however, if Source'First is not in Positive then Constraint_Error is raised.	68/1
function Translate (Source : in String; Mapping : in Maps.Character_Mapping)	69
return String;	
function Translate (Source : in String; Mapping : in Maps.Character_Mapping_Function)	69
return String;	
Returns the string S whose length is Source'Length and such that S(I) is the character to which Mapping maps the corresponding element of Source, for I in 1..Source'Length.	70
procedure Translate (Source : in out String; Mapping : in Maps.Character_Mapping);	71
procedure Translate (Source : in out String; Mapping : in Maps.Character_Mapping_Function);	71
Equivalent to Source := Translate(Source, Mapping).	72
function Replace_Slice (Source : in String; Low : in Positive; High : in Natural; By : in String)	73
return String;	
If Low > Source'Last+1, or High < Source'First-1, then Index_Error is propagated. Otherwise:	74/1
<ul style="list-style-type: none"> • If High >= Low, then the returned string comprises Source(Source'First..Low-1) & By & Source(High+1..Source'Last), but with lower bound 1. • If High < Low, then the returned string is Insert(Source, Before=>Low, New Item=>By). 	74.1/1

```

75      procedure Replace_Slice (Source    : in out String;
                                Low       : in Positive;
                                High      : in Natural;
                                By        : in String;
                                Drop      : in Truncation := Error;
                                Justify   : in Alignment := Left;
                                Pad       : in Character := Space);

76      Equivalent to Move(Replace_Slice(Source, Low, High, By), Source, Drop, Justify, Pad).

77      function Insert  (Source    : in String;
                        Before    : in Positive;
                        New_Item : in String)
                    return String;

78      Propagates Index_Error if Before is not in Source'First .. Source'Last+1; otherwise returns
          Source(Source'First..Before-1) & New_Item & Source(Before..Source'Last), but with lower
          bound 1.

79      procedure Insert  (Source    : in out String;
                        Before    : in Positive;
                        New_Item : in String;
                        Drop      : in Truncation := Error);

80      Equivalent to Move(Insert(Source, Before, New_Item), Source, Drop).

81      function Overwrite (Source    : in String;
                            Position : in Positive;
                            New_Item : in String)
                        return String;

82      Propagates Index_Error if Position is not in Source'First .. Source'Last+1; otherwise returns the
          string obtained from Source by consecutively replacing characters starting at Position with
          corresponding characters from New_Item. If the end of Source is reached before the characters
          in New_Item are exhausted, the remaining characters from New_Item are appended to the string.

83      procedure Overwrite (Source    : in out String;
                            Position : in Positive;
                            New_Item : in String;
                            Drop      : in Truncation := Right);

84      Equivalent to Move(Overwrite(Source, Position, New_Item), Source, Drop).

85      function Delete  (Source    : in String;
                        From      : in Positive;
                        Through   : in Natural)
                    return String;

86/1     If From <= Through, the returned string is Replace_Slice(Source, From, Through, ""), otherwise
          it is Source with lower bound 1.

87      procedure Delete  (Source    : in out String;
                            From      : in Positive;
                            Through   : in Natural;
                            Justify   : in Alignment := Left;
                            Pad       : in Character := Space);

88      Equivalent to Move>Delete(Source, From, Through), Source, Justify => Justify, Pad => Pad).

```

function Trim (Source : in String; Side : in Trim_End) return String;	89
Returns the string obtained by removing from Source all leading Space characters (if Side = Left), all trailing Space characters (if Side = Right), or all leading and trailing Space characters (if Side = Both).	90
procedure Trim (Source : in out String; Side : in Trim_End; Justify : in Alignment := Left; Pad : in Character := Space);	91
Equivalent to Move(Trim(Source, Side), Source, Justify=>Justify, Pad=>Pad).	92
function Trim (Source : in String; Left : in Maps.Character_Set; Right : in Maps.Character_Set) return String;	93
Returns the string obtained by removing from Source all leading characters in Left and all trailing characters in Right.	94
procedure Trim (Source : in out String; Left : in Maps.Character_Set; Right : in Maps.Character_Set; Justify : in Alignment := Strings.Left; Pad : in Character := Space);	95
Equivalent to Move(Trim(Source, Left, Right), Source, Justify => Justify, Pad=>Pad).	96
function Head (Source : in String; Count : in Natural; Pad : in Character := Space) return String;	97
Returns a string of length Count. If Count <= Source'Length, the string comprises the first Count characters of Source. Otherwise its contents are Source concatenated with Count-Source'Length Pad characters.	98
procedure Head (Source : in out String; Count : in Natural; Justify : in Alignment := Left; Pad : in Character := Space);	99
Equivalent to Move(Head(Source, Count, Pad), Source, Drop=>Error, Justify=>Justify, Pad=>Pad).	100
function Tail (Source : in String; Count : in Natural; Pad : in Character := Space) return String;	101
Returns a string of length Count. If Count <= Source'Length, the string comprises the last Count characters of Source. Otherwise its contents are Count-Source'Length Pad characters concatenated with Source.	102
procedure Tail (Source : in out String; Count : in Natural; Justify : in Alignment := Left; Pad : in Character := Space);	103
Equivalent to Move(Tail(Source, Count, Pad), Source, Drop=>Error, Justify=>Justify, Pad=>Pad).	104

```

105      function "*" (Left : in Natural;
                      Right : in Character) return String;
106      function "*" (Left : in Natural;
                      Right : in String) return String;

```

106/1 These functions replicate a character or string a specified number of times. The first function returns a string whose length is Left and each of whose elements is Right. The second function returns a string whose length is Left*Right'Length and whose value is the null string if Left = 0 and otherwise is (Left-1)*Right & Right with lower bound 1.

NOTES

- 107 9 In the Index and Count functions taking Pattern and Mapping parameters, the actual String parameter passed to Pattern should comprise characters occurring as target characters of the mapping. Otherwise the pattern will not match.
- 108 10 In the Insert subprograms, inserting at the end of a string is obtained by passing Source'Last+1 as the Before parameter.
- 109 11 If a null Character_Mapping_Function is passed to any of the string handling subprograms, Constraint_Error is propagated.

A.4.4 Bounded-Length String Handling

- 1 The language-defined package Strings.Bounded provides a generic package each of whose instances yields a private type Bounded_String and a set of operations. An object of a particular Bounded_String type represents a String whose low bound is 1 and whose length can vary conceptually between 0 and a maximum size established at the generic instantiation. The subprograms for fixed-length string handling are either overloaded directly for Bounded_String, or are modified as needed to reflect the variability in length. Additionally, since the Bounded_String type is private, appropriate constructor and selector operations are provided.

Static Semantics

- 2 The library package Strings.Bounded has the following declaration:

```

3      with Ada.Strings.Maps;
4      package Ada.Strings.Bounded is
5          pragma Preelaborate(Bounded);
6          generic
7              Max : Positive;      -- Maximum length of a Bounded_String
8          package Generic_Bounded_Length is
9              Max_Length : constant Positive := Max;
10             type Bounded_String is private;
11             Null_Bounded_String : constant Bounded_String;
12             subtype Length_Range is Natural range 0 .. Max_Length;
13             function Length (Source : in Bounded_String) return Length_Range;
14             -- Conversion, Concatenation, and Selection functions
15             function To_Bounded_String (Source : in String;
16                                         Drop    : in Truncation := Error)
17                 return Bounded_String;
18             function To_String (Source : in Bounded_String) return String;
19             procedure Set_Bounded_String
20                 (Target : out Bounded_String;
21                  Source : in String;
22                  Drop   : in Truncation := Error);

```

```

function Append (Left, Right : in Bounded_String;           13
                 Drop       : in Truncation := Error)
    return Bounded_String;
function Append (Left : in Bounded_String;                  14
                 Right : in String;
                 Drop  : in Truncation := Error)
    return Bounded_String;
function Append (Left : in String;                        15
                 Right : in Bounded_String;
                 Drop  : in Truncation := Error)
    return Bounded_String;
function Append (Left : in Bounded_String;                16
                 Right : in Character;
                 Drop  : in Truncation := Error)
    return Bounded_String;
function Append (Left : in Character;                     17
                 Right : in Bounded_String;
                 Drop  : in Truncation := Error)
    return Bounded_String;
procedure Append (Source : in out Bounded_String;          18
                  New_Item : in Bounded_String;
                  Drop    : in Truncation := Error);
procedure Append (Source : in out Bounded_String;          19
                  New_Item : in String;
                  Drop    : in Truncation := Error);
procedure Append (Source : in out Bounded_String;          20
                  New_Item : in Character;
                  Drop    : in Truncation := Error);
function "&" (Left, Right : in Bounded_String)            21
    return Bounded_String;
function "&" (Left : in Bounded_String; Right : in String) 22
    return Bounded_String;
function "&" (Left : in String; Right : in Bounded_String) 23
    return Bounded_String;
function "&" (Left : in Bounded_String; Right : in Character) 24
    return Bounded_String;
function "&" (Left : in Character; Right : in Bounded_String) 25
    return Bounded_String;
function Element (Source : in Bounded_String;              26
                  Index   : in Positive)
    return Character;
procedure Replace_Element (Source : in out Bounded_String; 27
                           Index   : in Positive;
                           By      : in Character);
function Slice (Source : in Bounded_String;                28
               Low    : in Positive;
               High   : in Natural)
    return String;
function Bounded_Slice                                     28.1/2
  (Source : in Bounded_String;
   Low    : in Positive;
   High   : in Natural)
    return Bounded_String;
procedure Bounded_Slice                                    28.2/2
  (Source : in Bounded_String;
   Target : out Bounded_String;
   Low    : in Positive;
   High   : in Natural);

```

```

29      function "="  (Left, Right : in Bounded_String) return Boolean;
function "="  (Left : in Bounded_String; Right : in String)
            return Boolean;
30      function "<"  (Left : in String; Right : in Bounded_String)
            return Boolean;
31      function "<"  (Left, Right : in Bounded_String) return Boolean;
32      function "<"  (Left : in Bounded_String; Right : in String)
            return Boolean;
33      function "<"  (Left : in String; Right : in Bounded_String)
            return Boolean;
34      function "<=" (Left, Right : in Bounded_String) return Boolean;
35      function "<=" (Left : in Bounded_String; Right : in String)
            return Boolean;
36      function "<=" (Left : in String; Right : in Bounded_String)
            return Boolean;
37      function ">"  (Left, Right : in Bounded_String) return Boolean;
38      function ">"  (Left : in Bounded_String; Right : in String)
            return Boolean;
39      function ">"  (Left : in String; Right : in Bounded_String)
            return Boolean;
40      function ">=" (Left, Right : in Bounded_String) return Boolean;
41      function ">=" (Left : in Bounded_String; Right : in String)
            return Boolean;
42      function ">=" (Left : in String; Right : in Bounded_String)
            return Boolean;
43/2    -- Search subprograms
43.1/2   function Index (Source  : in Bounded_String;
                           Pattern : in String;
                           From    : in Positive;
                           Going   : in Direction := Forward;
                           Mapping  : in Maps.Character_Mapping := Maps.Identity)
                  return Natural;
43.2/2   function Index (Source  : in Bounded_String;
                           Pattern : in String;
                           From    : in Positive;
                           Going   : in Direction := Forward;
                           Mapping  : in Maps.Character_Mapping_Function)
                  return Natural;
44       function Index (Source  : in Bounded_String;
                           Pattern : in String;
                           Going   : in Direction := Forward;
                           Mapping  : in Maps.Character_Mapping
                                     := Maps.Identity)
                  return Natural;
45       function Index (Source  : in Bounded_String;
                           Pattern : in String;
                           Going   : in Direction := Forward;
                           Mapping  : in Maps.Character_Mapping_Function)
                  return Natural;
45.1/2   function Index (Source  : in Bounded_String;
                           Set     : in Maps.Character_Set;
                           From   : in Positive;
                           Test   : in Membership := Inside;
                           Going  : in Direction := Forward)
                  return Natural;

```

```

function Index (Source : in Bounded_String;
                Set   : in Maps.Character_Set;
                Test  : in Membership := Inside;
                Going : in Direction  := Forward)
return Natural;                                         46

function Index_Non_Bank (Source : in Bounded_String;
                           From   : in Positive;
                           Going  : in Direction := Forward)
return Natural;                                         46.1/2

function Index_Non_Bank (Source : in Bounded_String;
                           Going  : in Direction := Forward)
return Natural;                                         47

function Count (Source   : in Bounded_String;
                 Pattern  : in String;
                 Mapping  : in Maps.Character_Mapping
                           := Maps.Identity)
return Natural;                                         48

function Count (Source   : in Bounded_String;
                 Pattern  : in String;
                 Mapping  : in Maps.Character_Mapping_Function)
return Natural;                                         49

function Count (Source   : in Bounded_String;
                 Set      : in Maps.Character_Set)
return Natural;                                         50

procedure Find_Token (Source : in Bounded_String;
                       Set   : in Maps.Character_Set;
                       Test  : in Membership;
                       First : out Positive;
                       Last  : out Natural);                                         51

-- String translation subprograms
function Translate (Source   : in Bounded_String;
                     Mapping  : in Maps.Character_Mapping)
return Bounded_String;                                         52
procedure Translate (Source   : in out Bounded_String;
                     Mapping  : in Maps.Character_Mapping);                                         53
function Translate (Source   : in Bounded_String;
                     Mapping  : in Maps.Character_Mapping_Function)
return Bounded_String;                                         54
procedure Translate (Source   : in out Bounded_String;
                     Mapping  : in Maps.Character_Mapping_Function);                                         55
-- String transformation subprograms
function Replace_Slice (Source   : in Bounded_String;
                         Low      : in Positive;
                         High     : in Natural;
                         By       : in String;
                         Drop    : in Truncation := Error)
return Bounded_String;                                         56
procedure Replace_Slice (Source   : in out Bounded_String;
                         Low      : in Positive;
                         High     : in Natural;
                         By       : in String;
                         Drop    : in Truncation := Error);                                         57
function Insert (Source   : in Bounded_String;
                  Before   : in Positive;
                  New_Item : in String;
                  Drop    : in Truncation := Error)
return Bounded_String;                                         58

```

```

61      procedure Insert (Source    : in out Bounded_String;
62                          Before     : in Positive;
63                          New_Item   : in String;
64                          Drop       : in Truncation := Error);
65
66      function Overwrite (Source    : in Bounded_String;
67                           Position   : in Positive;
68                           New_Item   : in String;
69                           Drop       : in Truncation := Error)
70                      return Bounded_String;
71
72      procedure Overwrite (Source    : in out Bounded_String;
73                           Position   : in Positive;
74                           New_Item   : in String;
75                           Drop       : in Truncation := Error);
76
77      function Delete (Source   : in Bounded_String;
78                       From      : in Positive;
79                       Through   : in Natural)
80                      return Bounded_String;
81
82      procedure Delete (Source   : in out Bounded_String;
83                       From      : in Positive;
84                       Through   : in Natural);
85
86      --String selector subprograms
87
88      function Trim (Source : in Bounded_String;
89                      Side   : in Trim_End)
90                      return Bounded_String;
91      procedure Trim (Source : in out Bounded_String;
92                      Side   : in Trim_End);
93
94      function Trim (Source : in Bounded_String;
95                      Left   : in Maps.Character_Set;
96                      Right  : in Maps.Character_Set)
97                      return Bounded_String;
98
99      procedure Trim (Source : in out Bounded_String;
100                     Left   : in Maps.Character_Set;
101                     Right  : in Maps.Character_Set);
102
103      function Head (Source : in Bounded_String;
104                      Count  : in Natural;
105                      Pad    : in Character  := Space;
106                      Drop   : in Truncation := Error)
107                      return Bounded_String;
108
109      procedure Head (Source : in out Bounded_String;
110                     Count  : in Natural;
111                     Pad    : in Character  := Space;
112                     Drop   : in Truncation := Error);
113
114      function Tail (Source : in Bounded_String;
115                      Count  : in Natural;
116                      Pad    : in Character  := Space;
117                      Drop   : in Truncation := Error)
118                      return Bounded_String;
119
120      procedure Tail (Source : in out Bounded_String;
121                     Count  : in Natural;
122                     Pad    : in Character  := Space;
123                     Drop   : in Truncation := Error);
124
125      --String constructor subprograms
126
127      function "*" (Left  : in Natural;
128                    Right : in Character)
129                    return Bounded_String;
130
131      function "*" (Left  : in Natural;
132                    Right : in String)
133                    return Bounded_String;

```

```

function "*" (Left  : in Natural;
              Right : in Bounded_String)
    return Bounded_String; 77

function Replicate (Count : in Natural;
                     Item  : in Character;
                     Drop   : in Truncation := Error)
    return Bounded_String; 78

function Replicate (Count : in Natural;
                     Item  : in String;
                     Drop   : in Truncation := Error)
    return Bounded_String; 79

function Replicate (Count : in Natural;
                     Item  : in Bounded_String;
                     Drop   : in Truncation := Error)
    return Bounded_String; 80

private
  ... -- not specified by the language
end Generic_Bounded_Length; 81

end Ada.Strings.Bounded; 82

```

Null_Bounded_String represents the null string. If an object of type Bounded_String is not otherwise initialized, it will be initialized to the same value as Null_Bounded_String. 83

```
function Length (Source : in Bounded_String) return Length_Range; 84
The Length function returns the length of the string represented by Source. 85
```

```
function To_Bounded_String (Source : in String;
                            Drop   : in Truncation := Error)
    return Bounded_String; 86
```

If Source'Length <= Max_Length then this function returns a Bounded_String that represents Source. Otherwise the effect depends on the value of Drop: 87

- If Drop=Left, then the result is a Bounded_String that represents the string comprising the rightmost Max_Length characters of Source. 88
- If Drop=Right, then the result is a Bounded_String that represents the string comprising the leftmost Max_Length characters of Source. 89
- If Drop=Error, then Strings.Length_Error is propagated. 90

```
function To_String (Source : in Bounded_String) return String; 91
```

To_String returns the String value with lower bound 1 represented by Source. If B is a Bounded_String, then B = To_Bounded_String(To_String(B)). 92

```
procedure Set_Bounded_String
  (Target : out Bounded_String;
   Source : in String;
   Drop   : in Truncation := Error); 92.1/2
```

Equivalent to Target := To_Bounded_String(Source, Drop); 92.2/2

Each of the Append functions returns a Bounded_String obtained by concatenating the string or character given or represented by one of the parameters, with the string or character given or represented by the other parameter, and applying To_Bounded_String to the concatenation result string, with Drop as provided to the Append function. 93

Each of the procedures Append(Source, New_Item, Drop) has the same effect as the corresponding assignment Source := Append(Source, New_Item, Drop). 94

- 95 Each of the "&" functions has the same effect as the corresponding Append function, with Error as the Drop parameter.

```
96  function Element (Source : in Bounded_String;
                     Index : in Positive)
            return Character;
```

97 Returns the character at position Index in the string represented by Source; propagates Index_Error if Index > Length(Source).

```
98  procedure Replace_Element (Source : in out Bounded_String;
                                Index : in Positive;
                                By     : in Character);
```

- 99 Updates Source such that the character at position Index in the string represented by Source is By; propagates Index_Error if Index > Length(Source).

```
100 function Slice (Source : in Bounded_String;
                      Low    : in Positive;
                      High   : in Natural)
                return String;
```

- 101/1 Returns the slice at positions Low through High in the string represented by Source; propagates Index_Error if Low > Length(Source)+1 or High > Length(Source). The bounds of the returned string are Low and High..

```
101.1/2 function Bounded_Slice
              (Source : in Bounded_String;
               Low    : in Positive;
               High   : in Natural)
                  return Bounded_String;
```

- 101.2/2 Returns the slice at positions Low through High in the string represented by Source as a bounded string; propagates Index_Error if Low > Length(Source)+1 or High > Length(Source).

```
101.3/2 procedure Bounded_Slice
              (Source : in Bounded_String;
               Target : out Bounded_String;
               Low    : in Positive;
               High   : in Natural);
```

- 101.4/2 Equivalent to Target := Bounded_Slice (Source, Low, High);

- 102 Each of the functions "=", "<", ">", "<=", and ">=" returns the same result as the corresponding String operation applied to the String values given or represented by the two parameters.

- 103 Each of the search subprograms (Index, Index_Non_Blank, Count, Find_Token) has the same effect as the corresponding subprogram in Strings.Fixed applied to the string represented by the Bounded_String parameter.

- 104 Each of the Translate subprograms, when applied to a Bounded_String, has an analogous effect to the corresponding subprogram in Strings.Fixed. For the Translate function, the translation is applied to the string represented by the Bounded_String parameter, and the result is converted (via To_Bounded_String) to a Bounded_String. For the Translate procedure, the string represented by the Bounded_String parameter after the translation is given by the Translate function for fixed-length strings applied to the string represented by the original value of the parameter.

- 105/1 Each of the transformation subprograms (Replace_Slice, Insert, Overwrite, Delete), selector subprograms (Trim, Head, Tail), and constructor functions ("*") has an effect based on its corresponding subprogram in Strings.Fixed, and Replicate is based on Fixed."*". In the case of a function, the corresponding fixed-

length string subprogram is applied to the string represented by the `Bounded_String` parameter. `To_Bounded_String` is applied the result string, with `Drop` (or `Error` in the case of `Generic_Bounded_Length."*"`) determining the effect when the string length exceeds `Max_Length`. In the case of a procedure, the corresponding function in `Strings.Bounded.Generic_Bounded_Length` is applied, with the result assigned into the `Source` parameter.

Implementation Advice

Bounded string objects should not be implemented by implicit pointers and dynamic allocation.

106

A.4.5 Unbounded-Length String Handling

The language-defined package `Strings.Unbounded` provides a private type `Unbounded_String` and a set of operations. An object of type `Unbounded_String` represents a String whose low bound is 1 and whose length can vary conceptually between 0 and `Natural'Last`. The subprograms for fixed-length string handling are either overloaded directly for `Unbounded_String`, or are modified as needed to reflect the flexibility in length. Since the `Unbounded_String` type is private, relevant constructor and selector operations are provided.

Static Semantics

The library package `Strings.Unbounded` has the following declaration:

```

with Ada.Strings.Maps;
package Ada.Strings.Unbounded is
  pragma Preelaborate(Unbounded);
  type Unbounded_String is private;
  pragma Preelaborable_Initialization(Unbounded_String);
  Null_Unbounded_String : constant Unbounded_String;
  function Length (Source : in Unbounded_String) return Natural;
  type String_Access is access all String;
  procedure Free (X : in out String_Access);

-- Conversion, Concatenation, and Selection functions
  function To_Unbounded_String (Source : in String)
    return Unbounded_String;
  function To_Unbounded_String (Length : in Natural)
    return Unbounded_String;
  function To_String (Source : in Unbounded_String) return String;
  procedure Set_Unbounded_String
    (Target : out Unbounded_String;
     Source : in String);
  procedure Append (Source   : in out Unbounded_String;
                  New_Item : in Unbounded_String);
  procedure Append (Source   : in out Unbounded_String;
                  New_Item : in String);
  procedure Append (Source   : in out Unbounded_String;
                  New_Item : in Character);
  function "&" (Left, Right : in Unbounded_String)
    return Unbounded_String;
  function "&" (Left : in Unbounded_String; Right : in String)
    return Unbounded_String;
  function "&" (Left : in String; Right : in Unbounded_String)
    return Unbounded_String;

```

```

18      function "&" (Left : in Unbounded_String; Right : in Character)
19          return Unbounded_String;
20      function "&" (Left : in Character; Right : in Unbounded_String)
21          return Unbounded_String;
22      function Element (Source : in Unbounded_String;
23                         Index   : in Positive)
24          return Character;
25      procedure Replace_Element (Source : in out Unbounded_String;
26                                  Index   : in Positive;
27                                  By      : in Character);
28      function Slice (Source : in Unbounded_String;
29                      Low    : in Positive;
30                      High   : in Natural)
31          return String;
32.1/2   function Unbounded_Slice
33         (Source : in Unbounded_String;
34          Low    : in Positive;
35          High   : in Natural)
36          return Unbounded_String;
37.2/2   procedure Unbounded_Slice
38         (Source : in     Unbounded_String;
39          Target : out    Unbounded_String;
40          Low    : in     Positive;
41          High   : in     Natural);
42
43      function "="  (Left, Right : in Unbounded_String) return Boolean;
44      function "="  (Left : in Unbounded_String; Right : in String)
45          return Boolean;
46      function "="  (Left : in String; Right : in Unbounded_String)
47          return Boolean;
48      function "<"  (Left, Right : in Unbounded_String) return Boolean;
49      function "<"  (Left : in Unbounded_String; Right : in String)
50          return Boolean;
51      function "<"  (Left : in String; Right : in Unbounded_String)
52          return Boolean;
53      function "<=" (Left, Right : in Unbounded_String) return Boolean;
54      function "<=" (Left : in Unbounded_String; Right : in String)
55          return Boolean;
56      function "<=" (Left : in String; Right : in Unbounded_String)
57          return Boolean;
58      function ">"  (Left, Right : in Unbounded_String) return Boolean;
59      function ">"  (Left : in Unbounded_String; Right : in String)
60          return Boolean;
61      function ">"  (Left : in String; Right : in Unbounded_String)
62          return Boolean;
63      function ">=" (Left, Right : in Unbounded_String) return Boolean;
64      function ">=" (Left : in Unbounded_String; Right : in String)
65          return Boolean;
66      function ">=" (Left : in String; Right : in Unbounded_String)
67          return Boolean;
68
69      -- Search subprograms
70.1/2   function Index (Source : in Unbounded_String;
71                        Pattern : in String;
72                        From    : in Positive;
73                        Going   : in Direction := Forward;
74                        Mapping  : in Maps.Character_Mapping := Maps.Identity)
75          return Natural;

```

```

function Index (Source : in Unbounded_String;
                Pattern : in String;
                From    : in Positive;
                Going   : in Direction := Forward;
                Mapping  : in Maps.Character_Mapping_Function) 38.2/2
  return Natural;
function Index (Source : in Unbounded_String;
                Pattern : in String;
                Going   : in Direction := Forward;
                Mapping  : in Maps.Character_Mapping
                           := Maps.Identity) 39
  return Natural;
function Index (Source : in Unbounded_String;
                Pattern : in String;
                Going   : in Direction := Forward;
                Mapping  : in Maps.Character_Mapping_Function) 40
  return Natural;
function Index (Source : in Unbounded_String;
                Set      : in Maps.Character_Set;
                From    : in Positive;
                Test    : in Membership := Inside;
                Going   : in Direction := Forward) 40.1/2
  return Natural;
function Index (Source : in Unbounded_String;
                Set      : in Maps.Character_Set;
                Test    : in Membership := Inside;
                Going   : in Direction := Forward) 41
  return Natural;
function Index_Non_Bank (Source : in Unbounded_String;
                           From   : in Positive;
                           Going  : in Direction := Forward) 41.1/2
  return Natural;
function Index_Non_Bank (Source : in Unbounded_String;
                           Going  : in Direction := Forward) 42
  return Natural;
function Count (Source : in Unbounded_String;
                  Pattern : in String;
                  Mapping  : in Maps.Character_Mapping
                           := Maps.Identity) 43
  return Natural;
function Count (Source : in Unbounded_String;
                  Pattern : in String;
                  Mapping  : in Maps.Character_Mapping_Function) 44
  return Natural;
function Count (Source : in Unbounded_String;
                  Set     : in Maps.Character_Set)
  return Natural; 45
procedure Find_Token (Source : in Unbounded_String;
                      Set     : in Maps.Character_Set;
                      Test    : in Membership;
                      First   : out Positive;
                      Last    : out Natural); 46
-- String translation subprograms 47
function Translate (Source : in Unbounded_String;
                     Mapping : in Maps.Character_Mapping) 48
  return Unbounded_String;
procedure Translate (Source : in out Unbounded_String;
                     Mapping : in Maps.Character_Mapping); 49
function Translate (Source : in Unbounded_String;
                     Mapping : in Maps.Character_Mapping_Function) 50
  return Unbounded_String;

```

```

51      procedure Translate (Source : in out Unbounded_String;
                           Mapping : in Maps.Character_Mapping_Function);
52      -- String transformation subprograms
53      function Replace_Slice (Source    : in Unbounded_String;
                               Low       : in Positive;
                               High      : in Natural;
                               By        : in String)
54          return Unbounded_String;
55      procedure Replace_Slice (Source    : in out Unbounded_String;
                               Low       : in Positive;
                               High      : in Natural;
                               By        : in String);
56      function Insert   (Source    : in Unbounded_String;
                           Before   : in Positive;
                           New_Item : in String)
57          return Unbounded_String;
58      procedure Insert   (Source    : in out Unbounded_String;
                           Before   : in Positive;
                           New_Item : in String);
59      function Overwrite (Source    : in Unbounded_String;
                           Position  : in Positive;
                           New_Item : in String)
60          return Unbounded_String;
61      procedure Overwrite (Source    : in out Unbounded_String;
                           Position  : in Positive;
                           New_Item : in String);
62      function Delete   (Source    : in Unbounded_String;
                           From     : in Positive;
                           Through  : in Natural)
63          return Unbounded_String;
64      procedure Delete   (Source    : in out Unbounded_String;
                           From     : in Positive;
                           Through  : in Natural);
65      function Trim     (Source    : in Unbounded_String;
                           Side     : in Trim_End)
66          return Unbounded_String;
67      procedure Trim     (Source    : in out Unbounded_String;
                           Side     : in Trim_End);
68      function Trim     (Source    : in Unbounded_String;
                           Left     : in Maps.Character_Set;
                           Right    : in Maps.Character_Set)
69          return Unbounded_String;
70      procedure Trim     (Source    : in out Unbounded_String;
                           Left     : in Maps.Character_Set;
                           Right    : in Maps.Character_Set);
71      function Head     (Source   : in Unbounded_String;
                           Count   : in Natural;
                           Pad     : in Character := Space)
72          return Unbounded_String;
73      procedure Head     (Source   : in out Unbounded_String;
                           Count   : in Natural;
                           Pad     : in Character := Space);
74      function Tail     (Source   : in Unbounded_String;
                           Count   : in Natural;
                           Pad     : in Character := Space)
75          return Unbounded_String;
76      procedure Tail     (Source   : in out Unbounded_String;
                           Count   : in Natural;
                           Pad     : in Character := Space);

```

```

function "*" (Left  : in Natural;
              Right : in Character)
    return Unbounded_String;
function "*" (Left  : in Natural;
              Right : in String)
    return Unbounded_String;
function "*" (Left  : in Natural;
              Right : in Unbounded_String)
    return Unbounded_String;
private
    ... -- not specified by the language
end Ada.Strings.Unbounded;

```

The type `Unbounded_String` needs finalization (see 7.6).

`Null_Unbounded_String` represents the null String. If an object of type `Unbounded_String` is not otherwise initialized, it will be initialized to the same value as `Null_Unbounded_String`.

The function `Length` returns the length of the String represented by `Source`.

The type `String_Access` provides a (non-private) access type for explicit processing of unbounded-length strings. The procedure `Free` performs an unchecked deallocation of an object of type `String_Access`.

The function `To_Unbounded_String`(`Source` : in `String`) returns an `Unbounded_String` that represents `Source`. The function `To_Unbounded_String`(`Length` : in `Natural`) returns an `Unbounded_String` that represents an uninitialized String whose length is `Length`.

The function `To_String` returns the String with lower bound 1 represented by `Source`. `To_String` and `To_Unbounded_String` are related as follows:

- If `S` is a `String`, then `To_String`(`To_Unbounded_String`(`S`)) = `S`.
- If `U` is an `Unbounded_String`, then `To_Unbounded_String`(`To_String`(`U`)) = `U`.

The procedure `Set_Unbounded_String` sets `Target` to an `Unbounded_String` that represents `Source`.

For each of the `Append` procedures, the resulting string represented by the `Source` parameter is given by the concatenation of the original value of `Source` and the value of `New_Item`.

Each of the "&" functions returns an `Unbounded_String` obtained by concatenating the string or character given or represented by one of the parameters, with the string or character given or represented by the other parameter, and applying `To_Unbounded_String` to the concatenation result string.

The `Element`, `Replace_Element`, and `Slice` subprograms have the same effect as the corresponding bounded-length string subprograms.

The function `Unbounded_Slice` returns the slice at positions `Low` through `High` in the string represented by `Source` as an `Unbounded_String`. The procedure `Unbounded_Slice` sets `Target` to the `Unbounded_String` representing the slice at positions `Low` through `High` in the string represented by `Source`. Both routines propagate `Index_Error` if `Low > Length(Source)+1` or `High > Length(Source)`.

Each of the functions "`=`", "`<`", "`>`", "`<=`", and "`>=`" returns the same result as the corresponding `String` operation applied to the `String` values given or represented by `Left` and `Right`.

Each of the search subprograms (`Index`, `Index_Non_Bank`, `Count`, `Find_Token`) has the same effect as the corresponding subprogram in `Strings.Fixed` applied to the string represented by the `Unbounded_String` parameter.

- 85 The Translate function has an analogous effect to the corresponding subprogram in Strings.Fixed. The translation is applied to the string represented by the Unbounded_String parameter, and the result is converted (via To_Unbounded_String) to an Unbounded_String.
- 86 Each of the transformation functions (Replace_Slice, Insert, Overwrite, Delete), selector functions (Trim, Head, Tail), and constructor functions ("*") is likewise analogous to its corresponding subprogram in Strings.Fixed. For each of the subprograms, the corresponding fixed-length string subprogram is applied to the string represented by the Unbounded_String parameter, and To_Unbounded_String is applied to the result string.
- 87 For each of the procedures Translate, Replace_Slice, Insert, Overwrite, Delete, Trim, Head, and Tail, the resulting string represented by the Source parameter is given by the corresponding function for fixed-length strings applied to the string represented by Source's original value.

Implementation Requirements

- 88 No storage associated with an Unbounded_String object shall be lost upon assignment or scope exit.

A.4.6 String-Handling Sets and Mappings

- 1 The language-defined package Strings.Maps.Constants declares Character_Set and Character_Mapping constants corresponding to classification and conversion functions in package Characters.Handling.

Static Semantics

- 2 The library package Strings.Maps.Constants has the following declaration:

```

3/2  package Ada.Strings.Maps.Constants is
      pragma Pure(Constants);
4
      Control_Set          : constant Character_Set;
      Graphic_Set          : constant Character_Set;
      Letter_Set            : constant Character_Set;
      Lower_Set             : constant Character_Set;
      Upper_Set              : constant Character_Set;
      Basic_Set              : constant Character_Set;
      Decimal_Digit_Set    : constant Character_Set;
      Hexadecimal_Digit_Set: constant Character_Set;
      Alphanumeric_Set      : constant Character_Set;
      Special_Set            : constant Character_Set;
      ISO_646_Set           : constant Character_Set;
5
      Lower_Case_Map        : constant Character_Mapping;
      --Maps to lower case for letters, else identity
      Upper_Case_Map        : constant Character_Mapping;
      --Maps to upper case for letters, else identity
      Basic_Map              : constant Character_Mapping;
      --Maps to basic letter for letters, else identity
6
      private
        ... -- not specified by the language
    end Ada.Strings.Maps.Constants;
```

- 7 Each of these constants represents a correspondingly named set of characters or character mapping in Characters.Handling (see A.3.2).

A.4.7 Wide_String Handling

- 1/2 Facilities for handling strings of Wide_Character elements are found in the packages Strings.Wide_Maps, Strings.Wide_Fixed, Strings.Wide_Bounded, Strings.Wide_Unbounded, and Strings.Wide_Maps.Wide_-

Constants, and in the functions `Strings.Wide_Hash`, `Strings.Wide_Fixed.Wide_Hash`, `Strings.Wide_Bounded.Wide_Hash`, and `Strings.Wide_Unbounded.Wide_Hash`. They provide the same string-handling operations as the corresponding packages and functions for strings of Character elements.

Static Semantics

The package `Strings.Wide_Maps` has the following declaration.

```

package Ada.Strings.Wide_Maps is
  pragma Preelaborate(Wide_Maps);
  -- Representation for a set of Wide_Character values:
  type Wide_Character_Set is private;
  pragma Preelaborable_Initialization(Wide_Character_Set);
  Null_Set : constant Wide_Character_Set;
  type Wide_Character_Range is
    record
      Low : Wide_Character;
      High : Wide_Character;
    end record;
  -- Represents Wide_Character range Low..High
  type Wide_Character_Ranges is array (Positive range <>)
    of Wide_Character_Range;
  function To_Set (Ranges : in Wide_Character_Ranges)
    return Wide_Character_Set;
  function To_Set (Span : in Wide_Character_Range)
    return Wide_Character_Set;
  function To_Ranges (Set : in Wide_Character_Set)
    return Wide_Character_Ranges;
  function "=" (Left, Right : in Wide_Character_Set) return Boolean;
  function "not" (Right : in Wide_Character_Set)
    return Wide_Character_Set;
  function "and" (Left, Right : in Wide_Character_Set)
    return Wide_Character_Set;
  function "or" (Left, Right : in Wide_Character_Set)
    return Wide_Character_Set;
  function "xor" (Left, Right : in Wide_Character_Set)
    return Wide_Character_Set;
  function "-" (Left, Right : in Wide_Character_Set)
    return Wide_Character_Set;
  function Is_In (Element : in Wide_Character;
                  Set : in Wide_Character_Set)
    return Boolean;
  function Is_Subset (Elements : in Wide_Character_Set;
                      Set : in Wide_Character_Set)
    return Boolean;
  function "<=" (Left : in Wide_Character_Set;
                Right : in Wide_Character_Set)
    return Boolean renames Is_Subset;
  -- Alternative representation for a set of Wide_Character values:
  subtype Wide_Character_Sequence is Wide_String;
  function To_Set (Sequence : in Wide_Character_Sequence)
    return Wide_Character_Set;
  function To_Set (Singleton : in Wide_Character)
    return Wide_Character_Set;
  function To_Sequence (Set : in Wide_Character_Set)
    return Wide_Character_Sequence;

```

```

20/2      -- Representation for a Wide_Character to Wide_Character mapping:
21       type Wide_Character_Mapping is private;
22       pragma Preelaborable_Initialization(Wide_Character_Mapping);
23       function Value (Map      : in Wide_Character_Mapping;
24                         Element : in Wide_Character)
25                     return Wide_Character;
26       Identity : constant Wide_Character_Mapping;
27       function To_Mapping (From, To : in Wide_Character_Sequence)
28                     return Wide_Character_Mapping;
29       function To_Domain (Map : in Wide_Character_Mapping)
30                     return Wide_Character_Sequence;
31       function To_Range  (Map : in Wide_Character_Mapping)
32                     return Wide_Character_Sequence;
33       type Wide_Character_Mapping_Function is
34           access function (From : in Wide_Character) return Wide_Character;
35       private
36           ... -- not specified by the language
37   end Ada.Strings.Wide_Maps;

```

28 The context clause for each of the packages Strings.Wide_Fixed, Strings.Wide_Bounded, and Strings.Wide_Unbounded identifies Strings.Wide_Maps instead of Strings.Maps.

29/2 For each of the packages Strings.Fixed, Strings.Bounded, Strings.Unbounded, and Strings.Maps.Constants, and for functions Strings.Hash, Strings.Fixed.Hash, Strings.Bounded.Hash, and Strings.Unbounded.Hash, the corresponding wide string package has the same contents except that

- 30 • Wide_Space replaces Space
- 31 • Wide_Character replaces Character
- 32 • Wide_String replaces String
- 33 • Wide_Character_Set replaces Character_Set
- 34 • Wide_Character_Mapping replaces Character_Mapping
- 35 • Wide_Character_Mapping_Function replaces Character_Mapping_Function
- 36 • Wide_Maps replaces Maps
- 37 • Bounded_Wide_String replaces Bounded_String
- 38 • Null_Bounded_Wide_String replaces Null_Bounded_String
- 39 • To_Bounded_Wide_String replaces To_Bounded_String
- 40 • To_Wide_String replaces To_String
- 40.1/2 • Set_Bounded_Wide_String replaces Set_Bounded_String
- 41 • Unbounded_Wide_String replaces Unbounded_String
- 42 • Null_Unbounded_Wide_String replaces Null_Unbounded_String
- 43 • Wide_String_Access replaces String_Access
- 44 • To_Unbounded_Wide_String replaces To_Unbounded_String
- 44.1/2 • Set_Unbounded_Wide_String replaces Set_Unbounded_String

45 The following additional declaration is present in Strings.Wide_Maps.Wide_Constants:

```

46/2      Character_Set : constant Wide_Maps.Wide_Character_Set;
47      -- Contains each Wide_Character value WC such that
48      -- Characters.Conversions.Is_Character(WC) is True

```

Each `Wide_Character_Set` constant in the package `Strings.Wide_Maps.Wide_Constants` contains no values outside the `Character` portion of `Wide_Character`. Similarly, each `Wide_Character_Mapping` constant in this package is the identity mapping when applied to any element outside the `Character` portion of `Wide_Character`.

Pragma Pure is replaced by pragma Preelaborate in `Strings.Wide_Maps.Wide_Constants`.

NOTES

12 If a null `Wide_Character_Mapping_Function` is passed to any of the `Wide_String` handling subprograms, `Constraint_Error` is propagated.

This paragraph was deleted.

A.4.8 Wide_Wide_String Handling

Facilities for handling strings of `Wide_Wide_Character` elements are found in the packages `Strings.Wide_Wide_Maps`, `Strings.Wide_Wide_Fixed`, `Strings.Wide_Wide_Bounded`, `Strings.Wide_Wide_Unbounded`, and `Strings.Wide_Wide_Maps.Wide_Wide_Constants`, and in the functions `Strings.Wide_Wide_Hash`, `Strings.Wide_Wide_Fixed.Wide_Wide_Hash`, `Strings.Wide_Wide_Bounded.Wide_Wide_Hash`, and `Strings.Wide_Wide_Unbounded.Wide_Wide_Hash`. They provide the same string-handling operations as the corresponding packages and functions for strings of `Character` elements.

Static Semantics

The library package `Strings.Wide_Wide_Maps` has the following declaration.

```
package Ada.Strings.Wide_Wide_Maps is
  pragma Preelaborate(Wide_Wide_Maps);
  -- Representation for a set of Wide_Wide_Character values:
  type Wide_Wide_Character_Set is private;
  pragma Preelaborable_Initialization(Wide_Wide_Character_Set);
  Null_Set : constant Wide_Wide_Character_Set;
  type Wide_Wide_Character_Range is
    record
      Low : Wide_Wide_Character;
      High : Wide_Wide_Character;
    end record;
  -- Represents Wide_Wide_Character range Low..High
  type Wide_Wide_Character_Ranges is array (Positive range <>)
    of Wide_Wide_Character_Range;
  function To_Set (Ranges : in Wide_Wide_Character_Ranges)
    return Wide_Wide_Character_Set;
  function To_Set (Span : in Wide_Wide_Character_Range)
    return Wide_Wide_Character_Set;
  function To_Ranges (Set : in Wide_Wide_Character_Set)
    return Wide_Wide_Character_Ranges;
  function "=" (Left, Right : in Wide_Wide_Character_Set) return Boolean;
  function "not" (Right : in Wide_Wide_Character_Set)
    return Wide_Wide_Character_Set;
  function "and" (Left, Right : in Wide_Wide_Character_Set)
    return Wide_Wide_Character_Set;
  function "or" (Left, Right : in Wide_Wide_Character_Set)
    return Wide_Wide_Character_Set;
  function "xor" (Left, Right : in Wide_Wide_Character_Set)
    return Wide_Wide_Character_Set;
  function "-" (Left, Right : in Wide_Wide_Character_Set)
    return Wide_Wide_Character_Set;
```

```

13/2      function Is_In (Element : in Wide_Wide_Character;
                      Set      : in Wide_Wide_Character_Set)
                     return Boolean;
14/2      function Is_Subset (Elements : in Wide_Wide_Character_Set;
                           Set      : in Wide_Wide_Character_Set)
                     return Boolean;
15/2      function "<=" (Left   : in Wide_Wide_Character_Set;
                      Right  : in Wide_Wide_Character_Set)
                     return Boolean renames Is_Subset;
16/2      -- Alternative representation for a set of Wide_Wide_Character values:
17/2      subtype Wide_Wide_Character_Sequence is Wide_Wide_String;
18/2      function To_Set (Sequence : in Wide_Wide_Character_Sequence)
                     return Wide_Wide_Character_Set;
19/2      function To_Set (Singleton : in Wide_Wide_Character)
                     return Wide_Wide_Character_Set;
20/2      function To_Sequence (Set : in Wide_Wide_Character_Set)
                     return Wide_Wide_Character_Sequence;
21/2      -- Representation for a Wide_Wide_Character to Wide_Wide_Character
22/2      -- mapping:
23/2      type Wide_Wide_Character_Mapping is private;
24/2      pragma Preelaborable_Initialization(Wide_Wide_Character_Mapping);
25/2      function Value (Map      : in Wide_Wide_Character_Mapping;
                      Element  : in Wide_Wide_Character)
                     return Wide_Wide_Character;
26/2      Identity : constant Wide_Wide_Character_Mapping;
27/2      function To_Mapping (From, To : in Wide_Wide_Character_Sequence)
                     return Wide_Wide_Character_Mapping;
28/2      function To_Domain (Map : in Wide_Wide_Character_Mapping)
                     return Wide_Wide_Character_Sequence;
29/2      function To_Range (Map : in Wide_Wide_Character_Mapping)
                     return Wide_Wide_Character_Sequence;
30/2      type Wide_Wide_Character_Mapping_Function is
31/2          access function (From : in Wide_Wide_Character)
                     return Wide_Wide_Character;
32/2      private
33/2          ... -- not specified by the language
34/2      end Ada.Strings.Wide_Wide_Maps;

```

28/2 The context clause for each of the packages `Strings.Wide_Wide_Fixed`, `Strings.Wide_Wide_Bounded`, and `Strings.Wide_Wide_Unbounded` identifies `Strings.Wide_Wide_Maps` instead of `StringsMaps`.

29/2 For each of the packages `Strings.Fixed`, `Strings.Bounded`, `Strings.Unbounded`, and `StringsMaps.Constants`, and for functions `Strings.Hash`, `Strings.Fixed.Hash`, `Strings.Bounded.Hash`, and `Strings.Unbounded.Hash`, the corresponding wide wide string package or function has the same contents except that

- 30/2 • `Wide_Wide_Space` replaces `Space`
- 31/2 • `Wide_Wide_Character` replaces `Character`
- 32/2 • `Wide_Wide_String` replaces `String`
- 33/2 • `Wide_Wide_Character_Set` replaces `Character_Set`
- 34/2 • `Wide_Wide_Character_Mapping` replaces `Character_Mapping`
- 35/2 • `Wide_Wide_Character_Mapping_Function` replaces `Character_Mapping_Function`
- 36/2 • `Wide_Wide_Maps` replaces `Maps`

• Bounded_Wide_Wide_String replaces Bounded_String	37/2
• Null_Bounded_Wide_Wide_String replaces Null_Bounded_String	38/2
• To_Bounded_Wide_Wide_String replaces To_Bounded_String	39/2
• To_Wide_Wide_String replaces To_String	40/2
• Set_Bounded_Wide_Wide_String replaces Set_Bounded_String	41/2
• Unbounded_Wide_Wide_String replaces Unbounded_String	42/2
• Null_Unbounded_Wide_Wide_String replaces Null_Unbounded_String	43/2
• Wide_Wide_String_Access replaces String_Access	44/2
• To_Unbounded_Wide_Wide_String replaces To_Unbounded_String	45/2
• Set_Unbounded_Wide_Wide_String replaces Set_Unbounded_String	46/2

The following additional declarations are present in Strings.Wide_Wide_Maps.Wide_Wide_Constants:

Character_Set : constant Wide_Wide_Maps.Wide_Wide_Character_Set;	48/2
-- Contains each Wide_Wide_Character value WWC such that	
-- Characters.Conversions.Is_Character(WWC) is True	
Wide_Character_Set : constant Wide_Wide_Maps.Wide_Wide_Character_Set;	
-- Contains each Wide_Wide_Character value WWC such that	
-- Characters.Conversions.Is_Wide_Character(WWC) is True	

Each Wide_Wide_Character_Set constant in the package Strings.Wide_Wide_Maps.Wide_Wide_Constants contains no values outside the Character portion of Wide_Wide_Character. Similarly, each Wide_Wide_Character_Mapping constant in this package is the identity mapping when applied to any element outside the Character portion of Wide_Wide_Character.

Pragma Pure is replaced by pragma Preelaborate in Strings.Wide_Wide_Maps.Wide_Wide_Constants.

NOTES

13 If a null Wide_Wide_Character_Mapping_Function is passed to any of the Wide_Wide_String handling subprograms, Constraint_Error is propagated.

A.4.9 String Hashing

Static Semantics

The library function Strings.Hash has the following declaration:

with Ada.Containers;	2/2
function Ada.Strings.Hash (Key : String) return Containers.Hash_Type;	
pragma Pure(Hash);	
>Returns an implementation-defined value which is a function of the value of Key. If A and B are strings such that A equals B, Hash(A) equals Hash(B).	3/2

The library function Strings.Fixed.Hash has the following declaration:

with Ada.Containers, Ada.Strings.Hash;	5/2
function Ada.Strings.Fixed.Hash (Key : String) return Containers.Hash_Type	
renames Ada.Strings.Hash;	
pragma Pure(Hash);	

6/2 The generic library function Strings.Bounded.Hash has the following declaration:

```
7/2   with Ada.Containers;
generic
  with package Bounded is
    new Ada.Strings.Bounded.Generic_Bounded_Length (<>);
  function Ada.Strings.Bounded.Hash (Key : Bounded.Bounded_String)
    return Containers.Hash_Type;
pragma Preelaborate(Hash);
```

8/2 Strings.Bounded.Hash is equivalent to the function call Strings.Hash (Bounded.To_String (Key));

9/2 The library function Strings.Unbounded.Hash has the following declaration:

```
10/2   with Ada.Containers;
  function Ada.Strings.Unbounded.Hash (Key : Unbounded_String)
    return Containers.Hash_Type;
pragma Preelaborate(Hash);
```

11/2 Strings.Unbounded.Hash is equivalent to the function call Strings.Hash (To_String (Key));

Implementation Advice

12/2 The Hash functions should be good hash functions, returning a wide spread of values for different string values. It should be unlikely for similar strings to return the same value.

A.5 The Numerics Packages

1 The library package Numerics is the parent of several child units that provide facilities for mathematical computation. One child, the generic package Generic_Elementary_Functions, is defined in A.5.1, together with nongeneric equivalents; two others, the package Float_Random and the generic package Discrete_Random, are defined in A.5.2. Additional (optional) children are defined in Annex G, “Numerics”.

Static Semantics

2/1 *This paragraph was deleted.*

```
3/2   package Ada.Numerics is
      pragma Pure(Numerics);
      Argument_Error : exception;
      Pi : constant :=
        3.14159_26535_89793_23846_26433_83279_50288_41971_69399_37511;
      π : constant := Pi;
      e : constant :=
        2.71828_18284_59045_23536_02874_71352_66249_77572_47093_69996;
    end Ada.Numerics;
```

4 The Argument_Error exception is raised by a subprogram in a child unit of Numerics to signal that one or more of the actual subprogram parameters are outside the domain of the corresponding mathematical function.

Implementation Permissions

5 The implementation may specify the values of Pi and e to a larger number of significant digits.

A.5.1 Elementary Functions

Implementation-defined approximations to the mathematical functions known as the “elementary functions” are provided by the subprograms in Numerics.Generic_Elementary_Functions. Nongeneric equivalents of this generic package for each of the predefined floating point types are also provided as children of Numerics.

Static Semantics

The generic library package Numerics.Generic_Elementary_Functions has the following declaration:

```

generic
  type Float_Type is digits <>;

package Ada.Numerics.Generic_Elementary_Functions is
  pragma Pure(Generic_Elementary_Functions);

  function Sqrt      (X          : Float_Type'Base) return Float_Type'Base;
  function Log       (X          : Float_Type'Base) return Float_Type'Base;
  function Log      (X, Base    : Float_Type'Base) return Float_Type'Base;
  function Exp       (X          : Float_Type'Base) return Float_Type'Base;
  function "***"   (Left, Right : Float_Type'Base) return Float_Type'Base;

  function Sin       (X          : Float_Type'Base) return Float_Type'Base;
  function Sin      (X, Cycle   : Float_Type'Base) return Float_Type'Base;
  function Cos       (X          : Float_Type'Base) return Float_Type'Base;
  function Cos      (X, Cycle   : Float_Type'Base) return Float_Type'Base;
  function Tan       (X          : Float_Type'Base) return Float_Type'Base;
  function Tan      (X, Cycle   : Float_Type'Base) return Float_Type'Base;
  function Cot       (X          : Float_Type'Base) return Float_Type'Base;
  function Cot      (X, Cycle   : Float_Type'Base) return Float_Type'Base;

  function Arcsin    (X          : Float_Type'Base) return Float_Type'Base;
  function Arcsin   (X, Cycle   : Float_Type'Base) return Float_Type'Base;
  function Arccos    (X          : Float_Type'Base) return Float_Type'Base;
  function Arccos   (X, Cycle   : Float_Type'Base) return Float_Type'Base;
  function Arctan    (Y          : Float_Type'Base;
                      X          : Float_Type'Base := 1.0) return Float_Type'Base;

  function Arctan   (Y          : Float_Type'Base;
                     X          : Float_Type'Base := 1.0;
                     Cycle     : Float_Type'Base) return Float_Type'Base;
  function Arccot    (X          : Float_Type'Base;
                      Y          : Float_Type'Base := 1.0) return Float_Type'Base;

  function Arccot   (X          : Float_Type'Base;
                     Y          : Float_Type'Base := 1.0;
                     Cycle     : Float_Type'Base) return Float_Type'Base;

  function Sinh      (X          : Float_Type'Base) return Float_Type'Base;
  function Cosh      (X          : Float_Type'Base) return Float_Type'Base;
  function Tanh      (X          : Float_Type'Base) return Float_Type'Base;
  function Coth      (X          : Float_Type'Base) return Float_Type'Base;
  function Arcsinh   (X          : Float_Type'Base) return Float_Type'Base;
  function Arccosh   (X          : Float_Type'Base) return Float_Type'Base;
  function Arctanh   (X          : Float_Type'Base) return Float_Type'Base;
  function Arccoth   (X          : Float_Type'Base) return Float_Type'Base;

end Ada.Numerics.Generic_Elementary_Functions;
```

The library package Numerics.Elementary_Functions is declared pure and defines the same subprograms as Numerics.Generic_Elementary_Functions, except that the predefined type `Float` is systematically substituted for `Float_Type'Base` throughout. Nongeneric equivalents of Numerics.Generic_Elementary_Functions for each of the other predefined floating point types are defined similarly, with the names `Numerics.Short_Elementary_Functions`, `Numerics.Long_Elementary_Functions`, etc.

- 10 The functions have their usual mathematical meanings. When the Base parameter is specified, the Log function computes the logarithm to the given base; otherwise, it computes the natural logarithm. When the Cycle parameter is specified, the parameter X of the forward trigonometric functions (Sin, Cos, Tan, and Cot) and the results of the inverse trigonometric functions (Arcsin, Arccos, Arctan, and Arccot) are measured in units such that a full cycle of revolution has the given value; otherwise, they are measured in radians.
- 11 The computed results of the mathematically multivalued functions are rendered single-valued by the following conventions, which are meant to imply the principal branch:
- 12 • The results of the Sqrt and Arccosh functions and that of the exponentiation operator are nonnegative.
 - 13 • The result of the Arcsin function is in the quadrant containing the point $(1.0, x)$, where x is the value of the parameter X. This quadrant is I or IV; thus, the range of the Arcsin function is approximately $-\pi/2.0$ to $\pi/2.0$ ($-Cycle/4.0$ to $Cycle/4.0$, if the parameter Cycle is specified).
 - 14 • The result of the Arccos function is in the quadrant containing the point $(x, 1.0)$, where x is the value of the parameter X. This quadrant is I or II; thus, the Arccos function ranges from 0.0 to approximately π ($Cycle/2.0$, if the parameter Cycle is specified).
 - 15 • The results of the Arctan and Arccot functions are in the quadrant containing the point (x, y) , where x and y are the values of the parameters X and Y, respectively. This may be any quadrant (I through IV) when the parameter X (resp., Y) of Arctan (resp., Arccot) is specified, but it is restricted to quadrants I and IV (resp., I and II) when that parameter is omitted. Thus, the range when that parameter is specified is approximately $-\pi$ to π ($-Cycle/2.0$ to $Cycle/2.0$, if the parameter Cycle is specified); when omitted, the range of Arctan (resp., Arccot) is that of Arcsin (resp., Arccos), as given above. When the point (x, y) lies on the negative x-axis, the result approximates
 - 16 • π (resp., $-\pi$) when the sign of the parameter Y is positive (resp., negative), if `Float_Type'Signed_Zeros` is True;
 - 17 • π , if `Float_Type'Signed_Zeros` is False.
- 18 (In the case of the inverse trigonometric functions, in which a result lying on or near one of the axes may not be exactly representable, the approximation inherent in computing the result may place it in an adjacent quadrant, close to but on the wrong side of the axis.)

Dynamic Semantics

- 19 The exception `Numerics.Argument_Error` is raised, signaling a parameter value outside the domain of the corresponding mathematical function, in the following cases:
- 20 • by any forward or inverse trigonometric function with specified cycle, when the value of the parameter Cycle is zero or negative;
 - 21 • by the Log function with specified base, when the value of the parameter Base is zero, one, or negative;
 - 22 • by the Sqrt and Log functions, when the value of the parameter X is negative;
 - 23 • by the exponentiation operator, when the value of the left operand is negative or when both operands have the value zero;
 - 24 • by the Arcsin, Arccos, and Arctanh functions, when the absolute value of the parameter X exceeds one;
 - 25 • by the Arctan and Arccot functions, when the parameters X and Y both have the value zero;
 - 26 • by the Arccosh function, when the value of the parameter X is less than one; and

- by the Arccoth function, when the absolute value of the parameter X is less than one. 27

The exception Constraint_Error is raised, signaling a pole of the mathematical function (analogous to dividing by zero), in the following cases, provided that Float_Type'Machine_Overflows is True:

- by the Log, Cot, and Coth functions, when the value of the parameter X is zero; 29
- by the exponentiation operator, when the value of the left operand is zero and the value of the exponent is negative; 30
- by the Tan function with specified cycle, when the value of the parameter X is an odd multiple of the quarter cycle; 31
- by the Cot function with specified cycle, when the value of the parameter X is zero or a multiple of the half cycle; and 32
- by the Arctanh and Arccoth functions, when the absolute value of the parameter X is one. 33

Constraint_Error can also be raised when a finite result overflows (see G.2.4); this may occur for parameter values sufficiently *near* poles, and, in the case of some of the functions, for parameter values with sufficiently large magnitudes. When Float_Type'Machine_Overflows is False, the result at poles is unspecified.

When one parameter of a function with multiple parameters represents a pole and another is outside the function's domain, the latter takes precedence (i.e., Numerics.Argument_Error is raised). 35

Implementation Requirements

In the implementation of Numerics.Generic_Elementary_Functions, the range of intermediate values allowed during the calculation of a final result shall not be affected by any range constraint of the subtype Float_Type. 36

In the following cases, evaluation of an elementary function shall yield the *prescribed result*, provided that the preceding rules do not call for an exception to be raised: 37

- When the parameter X has the value zero, the Sqrt, Sin, Arcsin, Tan, Sinh, Arcsinh, Tanh, and Arctanh functions yield a result of zero, and the Exp, Cos, and Cosh functions yield a result of one. 38
- When the parameter X has the value one, the Sqrt function yields a result of one, and the Log, Arccos, and Arccosh functions yield a result of zero. 39
- When the parameter Y has the value zero and the parameter X has a positive value, the Arctan and Arccot functions yield a result of zero. 40
- The results of the Sin, Cos, Tan, and Cot functions with specified cycle are exact when the mathematical result is zero; those of the first two are also exact when the mathematical result is ± 1.0 . 41
- Exponentiation by a zero exponent yields the value one. Exponentiation by a unit exponent yields the value of the left operand. Exponentiation of the value one yields the value one. Exponentiation of the value zero yields the value zero. 42

Other accuracy requirements for the elementary functions, which apply only in implementations conforming to the Numerics Annex, and then only in the “strict” mode defined there (see G.2), are given in G.2.4. 43

When Float_Type'Signed_Zeros is True, the sign of a zero result shall be as follows: 44

- 45 • A prescribed zero result delivered *at the origin* by one of the odd functions (Sin, Arcsin, Sinh, Arcsinh, Tan, Arctan or Arccot as a function of Y when X is fixed and positive, Tanh, and Arctanh) has the sign of the parameter X (Y, in the case of Arctan or Arccot).
- 46 • A prescribed zero result delivered by one of the odd functions *away from the origin*, or by some other elementary function, has an implementation-defined sign.
- 47 • A zero result that is not a prescribed result (i.e., one that results from rounding or underflow) has the correct mathematical sign.

Implementation Permissions

- 48 The nongeneric equivalent packages may, but need not, be actual instantiations of the generic package for the appropriate predefined type.

A.5.2 Random Number Generation

- 1 Facilities for the generation of pseudo-random floating point numbers are provided in the package Numerics.Float_Random; the generic package Numerics.Discrete_Random provides similar facilities for the generation of pseudo-random integers and pseudo-random values of enumeration types. For brevity, pseudo-random values of any of these types are called *random numbers*.
- 2 Some of the facilities provided are basic to all applications of random numbers. These include a limited private type each of whose objects serves as the generator of a (possibly distinct) sequence of random numbers; a function to obtain the “next” random number from a given sequence of random numbers (that is, from its generator); and subprograms to initialize or reinitialize a given generator to a time-dependent state or a state denoted by a single integer.
- 3 Other facilities are provided specifically for advanced applications. These include subprograms to save and restore the state of a given generator; a private type whose objects can be used to hold the saved state of a generator; and subprograms to obtain a string representation of a given generator state, or, given such a string representation, the corresponding state.

Static Semantics

- 4 The library package Numerics.Float_Random has the following declaration:

```

5   package Ada.Numerics.Float_Random is
6     -- Basic facilities
7     type Generator is limited private;
8     subtype Uniformly_Distributed is Float range 0.0 .. 1.0;
9     function Random (Gen : Generator) return Uniformly_Distributed;
10    procedure Reset (Gen      : in Generator;
11                      Initiator : in Integer);
12    procedure Reset (Gen      : in Generator);
13    -- Advanced facilities
14    type State is private;
15    procedure Save   (Gen      : in Generator;
16                      To_State : out State);
17    procedure Reset (Gen      : in Generator;
18                      From_State : in State);
19    Max_Image_Width : constant := implementation-defined integer value;
20    function Image  (Of_State  : State)  return String;
21    function Value   (Coded_State : String) return State;

```

```
private
  ... -- not specified by the language
end Ada.Numerics.Float_Random;
```

The type Generator needs finalization (see 7.6).

The generic library package Numerics.Discrete_Random has the following declaration:

```
generic
  type Result_Subtype is (<>);
package Ada.Numerics.Discrete_Random is
  -- Basic facilities
  type Generator is limited private;
  function Random (Gen : Generator) return Result_Subtype;
  procedure Reset (Gen      : in Generator;
                  Initiator : in Integer);
  procedure Reset (Gen      : in Generator);
  -- Advanced facilities
  type State is private;
  procedure Save   (Gen      : in Generator;
                  To_State  : out State);
  procedure Reset (Gen      : in Generator;
                  From_State : in State);
  Max_Image_Width : constant := implementation-defined integer value;
  function Image (Of_State  : State) return String;
  function Value (Coded_State : String) return State;
private
  ... -- not specified by the language
end Ada.Numerics.Discrete_Random;
```

The type Generator needs finalization (see 7.6) in every instantiation of Numerics.Discrete_Random.

An object of the limited private type Generator is associated with a sequence of random numbers. Each generator has a hidden (internal) state, which the operations on generators use to determine the position in the associated sequence. All generators are implicitly initialized to an unspecified state that does not vary from one program execution to another; they may also be explicitly initialized, or reinitialized, to a time-dependent state, to a previously saved state, or to a state uniquely denoted by an integer value.

An object of the private type State can be used to hold the internal state of a generator. Such objects are only needed if the application is designed to save and restore generator states or to examine or manufacture them.

The operations on generators affect the state and therefore the future values of the associated sequence. The semantics of the operations on generators and states are defined below.

```
function Random (Gen : Generator) return Uniformly_Distributed;
function Random (Gen : Generator) return Result_Subtype;
```

Obtains the “next” random number from the given generator, relative to its current state, according to an implementation-defined algorithm. The result of the function in Numerics.Float_Random is delivered as a value of the subtype Uniformly_Distributed, which is a subtype of the predefined type Float having a range of 0.0 .. 1.0. The result of the function in an instantiation of Numerics.Discrete_Random is delivered as a value of the generic formal subtype Result_Subtype.

33 **procedure** Reset (Gen : **in** Generator;
 Initiator : **in** Integer);
procedure Reset (Gen : **in** Generator);

34 Sets the state of the specified generator to one that is an unspecified function of the value of the parameter Initiator (or to a time-dependent state, if only a generator parameter is specified). The latter form of the procedure is known as the *time-dependent Reset procedure*.

35 **procedure** Save (Gen : **in** Generator;
 To_State : **out** State);
procedure Reset (Gen : **in** Generator;
 From_State : **in** State);

36 Save obtains the current state of a generator. Reset gives a generator the specified state. A generator that is reset to a state previously obtained by invoking Save is restored to the state it had when Save was invoked.

37 **function** Image (Of_State : State) **return** String;
function Value (Coded_State : String) **return** State;

38 Image provides a representation of a state coded (in an implementation-defined way) as a string whose length is bounded by the value of Max_Image_Width. Value is the inverse of Image: Value(Image(S)) = S for each state S that can be obtained from a generator by invoking Save.

Dynamic Semantics

39 Instantiation of Numerics.Discrete_Random with a subtype having a null range raises Constraint_Error.

40/1 *This paragraph was deleted.*

Bounded (Run-Time) Errors

40.1/1 It is a bounded error to invoke Value with a string that is not the image of any generator state. If the error is detected, Constraint_Error or Program_Error is raised. Otherwise, a call to Reset with the resulting state will produce a generator such that calls to Random with this generator will produce a sequence of values of the appropriate subtype, but which might not be random in character. That is, the sequence of values might not fulfill the implementation requirements of this subclause.

Implementation Requirements

41 A sufficiently long sequence of random numbers obtained by successive calls to Random is approximately uniformly distributed over the range of the result subtype.

42 The Random function in an instantiation of Numerics.Discrete_Random is guaranteed to yield each value in its result subtype in a finite number of calls, provided that the number of such values does not exceed 2¹⁵.

43 Other performance requirements for the random number generator, which apply only in implementations conforming to the Numerics Annex, and then only in the “strict” mode defined there (see G.2), are given in G.2.5.

Documentation Requirements

44 No one algorithm for random number generation is best for all applications. To enable the user to determine the suitability of the random number generators for the intended application, the implementation shall describe the algorithm used and shall give its period, if known exactly, or a lower bound on the period, if the exact period is unknown. Periods that are so long that the periodicity is unobservable in practice can be described in such terms, without giving a numerical bound.

The implementation also shall document the minimum time interval between calls to the time-dependent Reset procedure that are guaranteed to initiate different sequences, and it shall document the nature of the strings that Value will accept without raising Constraint_Error.

Implementation Advice

Any storage associated with an object of type Generator should be reclaimed on exit from the scope of the object.

If the generator period is sufficiently long in relation to the number of distinct initiator values, then each possible value of Initiator passed to Reset should initiate a sequence of random numbers that does not, in a practical sense, overlap the sequence initiated by any other value. If this is not possible, then the mapping between initiator values and generator states should be a rapidly varying function of the initiator value.

NOTES

14 If two or more tasks are to share the same generator, then the tasks have to synchronize their access to the generator as for any shared variable (see 9.10).

15 Within a given implementation, a repeatable random number sequence can be obtained by relying on the implicit initialization of generators or by explicitly initializing a generator with a repeatable initiator value. Different sequences of random numbers can be obtained from a given generator in different program executions by explicitly initializing the generator to a time-dependent state.

16 A given implementation of the Random function in Numerics.Float_Random may or may not be capable of delivering the values 0.0 or 1.0. Portable applications should assume that these values, or values sufficiently close to them to behave indistinguishably from them, can occur. If a sequence of random integers from some fixed range is needed, the application should use the Random function in an appropriate instantiation of Numerics.Discrete_Random, rather than transforming the result of the Random function in Numerics.Float_Random. However, some applications with unusual requirements, such as for a sequence of random integers each drawn from a different range, will find it more convenient to transform the result of the floating point Random function. For $M \geq 1$, the expression

`Integer(Float(M) * Random(G)) mod M`

transforms the result of Random(G) to an integer uniformly distributed over the range 0 .. $M-1$; it is valid even if Random delivers 0.0 or 1.0. Each value of the result range is possible, provided that M is not too large. Exponentially distributed (floating point) random numbers with mean and standard deviation 1.0 can be obtained by the transformation

`-Log(Random(G)) + Float'Model_Small`

where Log comes from Numerics.Elementary_Functions (see A.5.1); in this expression, the addition of Float'Model_Small avoids the exception that would be raised were Log to be given the value zero, without affecting the result (in most implementations) when Random returns a nonzero value.

Examples

Example of a program that plays a simulated dice game:

```
with Ada.Numerics.Discrete_Random;
procedure Dice_Game is
    subtype Die is Integer range 1 .. 6;
    subtype Dice is Integer range 2*Die'First .. 2*Die'Last;
    package Random_Die is new Ada.Numerics.Discrete_Random (Die);
    use Random_Die;
    G : Generator;
    D : Dice;
begin
    Reset (G); -- Start the generator in a unique state in each run
loop
    -- Roll a pair of dice; sum and process the results
    D := Random(G) + Random(G);
    . . .
end loop;
end Dice_Game;
```

57 Example of a program that simulates coin tosses:

```

58  with Ada.Numerics.Discrete_Random;
59  procedure Flip_A_Coin is
60    type Coin is (Heads, Tails);
61    package Random_Coin is new Ada.Numerics.Discrete_Random (Coin);
62    use Random_Coin;
63    G : Generator;
64
65 begin
66   Reset (G); -- Start the generator in a unique state in each run
67   loop
68     -- Toss a coin and process the result
69     case Random(G) is
70       when Heads =>
71         ...
72       when Tails =>
73         ...
74     end case;
75   end loop;
76 end Flip_A_Coin;
```

59 Example of a parallel simulation of a physical system, with a separate generator of event probabilities in each task:

```

60  with Ada.Numerics.Float_Random;
61  procedure Parallel_Simulation is
62    use Ada.Numerics.Float_Random;
63    task type Worker is
64      entry Initialize_Generator (Initiator : in Integer);
65      ...
66    end Worker;
67    W : array (1 .. 10) of Worker;
68    task body Worker is
69      G : Generator;
70      Probability_Of_Event : Uniformly_Distributed;
71
72 begin
73   accept Initialize_Generator (Initiator : in Integer) do
74     Reset (G, Initiator);
75   end Initialize_Generator;
76   loop
77     ...
78     Probability_Of_Event := Random(G);
79     ...
80   end loop;
81 end Worker;
82 begin
83   -- Initialize the generators in the Worker tasks to different states
84   for I in W'Range loop
85     W(I).Initialize_Generator (I);
86   end loop;
87   ... -- Wait for the Worker tasks to terminate
88 end Parallel_Simulation;
```

NOTES

61 17 Notes on the last example: Although each Worker task initializes its generator to a different state, those states will be the same in every execution of the program. The generator states can be initialized uniquely in each program execution by instantiating Ada.Numerics.Discrete_Random for the type Integer in the main procedure, resetting the generator obtained from that instance to a time-dependent state, and then using random integers obtained from that generator to initialize the generators in each Worker task.

A.5.3 Attributes of Floating Point Types

Static Semantics

The following *representation-oriented attributes* are defined for every subtype S of a floating point type T.

S'Machine_Radix

Yields the radix of the hardware representation of the type T. The value of this attribute is of the type *universal_integer*.

The values of other representation-oriented attributes of a floating point subtype, and of the “primitive function” attributes of a floating point subtype described later, are defined in terms of a particular representation of nonzero values called the *canonical form*. The canonical form (for the type T) is the form

$$\pm \text{mantissa} \cdot T\text{Machine_Radix}^{\text{exponent}}$$

where

- *mantissa* is a fraction in the number base *TMachine_Radix*, the first digit of which is nonzero,
- and
- *exponent* is an integer.

S'Machine_Mantissa

Yields the largest value of *p* such that every value expressible in the canonical form (for the type T), having a *p*-digit *mantissa* and an *exponent* between *TMachine_Emin* and *TMachine_Emax*, is a machine number (see 3.5.7) of the type T. This attribute yields a value of the type *universal_integer*.

S'Machine_Emin

Yields the smallest (most negative) value of *exponent* such that every value expressible in the canonical form (for the type T), having a *mantissa* of *TMachine_Mantissa* digits, is a machine number (see 3.5.7) of the type T. This attribute yields a value of the type *universal_integer*.

S'Machine_Emax

Yields the largest (most positive) value of *exponent* such that every value expressible in the canonical form (for the type T), having a *mantissa* of *TMachine_Mantissa* digits, is a machine number (see 3.5.7) of the type T. This attribute yields a value of the type *universal_integer*.

S'Denorm

Yields the value True if every value expressible in the form

$$\pm \text{mantissa} \cdot T\text{Machine_Radix}^{T\text{Machine_Emin}}$$

where *mantissa* is a nonzero *TMachine_Mantissa*-digit fraction in the number base *TMachine_Radix*, the first digit of which is zero, is a machine number (see 3.5.7) of the type T; yields the value False otherwise. The value of this attribute is of the predefined type Boolean.

The values described by the formula in the definition of S'Denorm are called *denormalized numbers*. A nonzero machine number that is not a denormalized number is a *normalized number*. A normalized number x of a given type T is said to be *represented in canonical form* when it is expressed in the canonical form (for the type T) with a *mantissa* having *TMachine_Mantissa* digits; the resulting form is the *canonical-form representation* of x.

S'Machine_Rounds

Yields the value True if rounding is performed on inexact results of every predefined operation that yields a result of the type T; yields the value False otherwise. The value of this attribute is of the predefined type Boolean.

- | | | |
|----|---------------------|--|
| 12 | S'Machine_Overflows | Yields the value True if overflow and divide-by-zero are detected and reported by raising Constraint_Error for every predefined operation that yields a result of the type T ; yields the value False otherwise. The value of this attribute is of the predefined type Boolean. |
| 13 | S'Signed_Zeros | Yields the value True if the hardware representation for the type T has the capability of representing both positively and negatively signed zeros, these being generated and used by the predefined operations of the type T as specified in IEC 559:1989; yields the value False otherwise. The value of this attribute is of the predefined type Boolean. |
| 14 | | For every value x of a floating point type T , the <i>normalized exponent</i> of x is defined as follows: |
| 15 | | <ul style="list-style-type: none"> • the normalized exponent of zero is (by convention) zero; |
| 16 | | <ul style="list-style-type: none"> • for nonzero x, the normalized exponent of x is the unique integer k such that $T\text{Machine_Radix}^{k-1} \leq x < T\text{Machine_Radix}^k$. |
| 17 | | The following <i>primitive function attributes</i> are defined for any subtype S of a floating point type T . |
| 18 | S'Exponent | S'Exponent denotes a function with the following specification: |
| 19 | | function S'Exponent ($X : T$)
return universal_integer |
| 20 | | The function yields the normalized exponent of X . |
| 21 | S'Fraction | S'Fraction denotes a function with the following specification: |
| 22 | | function S'Fraction ($X : T$)
return T |
| 23 | | The function yields the value $X \cdot T\text{Machine_Radix}^{-k}$, where k is the normalized exponent of X . A zero result, which can only occur when X is zero, has the sign of X . |
| 24 | S'Compose | S'Compose denotes a function with the following specification: |
| 25 | | function S'Compose ($Fraction : T$;
Exponent : universal_integer)
return T |
| 26 | | Let v be the value $Fraction \cdot T\text{Machine_Radix}^{Exponent-k}$, where k is the normalized exponent of $Fraction$. If v is a machine number of the type T , or if $ v \geq T\text{Model_Small}$, the function yields v ; otherwise, it yields either one of the machine numbers of the type T adjacent to v . Constraint_Error is optionally raised if v is outside the base range of S . A zero result has the sign of $Fraction$ when S'Signed_Zeros is True. |
| 27 | S'Scaling | S'Scaling denotes a function with the following specification: |
| 28 | | function S'Scaling ($X : T$;
Adjustment : universal_integer)
return T |
| 29 | | Let v be the value $X \cdot T\text{Machine_Radix}^{Adjustment}$. If v is a machine number of the type T , or if $ v \geq T\text{Model_Small}$, the function yields v ; otherwise, it yields either one of the machine numbers of the type T adjacent to v . Constraint_Error is optionally raised if v is outside the base range of S . A zero result has the sign of X when S'Signed_Zeros is True. |
| 30 | S'Floor | S'Floor denotes a function with the following specification: |
| 31 | | function S'Floor ($X : T$)
return T |
| 32 | | The function yields the value $\lfloor X \rfloor$, i.e., the largest (most positive) integral value less than or equal to X . When X is zero, the result has the sign of X ; a zero result otherwise has a positive sign. |

S'Ceiling	S'Ceiling denotes a function with the following specification:	33
	<pre>function S'Ceiling (X : T) return T</pre>	34
	The function yields the value $\lceil X \rceil$, i.e., the smallest (most negative) integral value greater than or equal to X . When X is zero, the result has the sign of X ; a zero result otherwise has a negative sign when S'Signed_Zeros is True.	35
S'Rounding	S'Rounding denotes a function with the following specification:	36
	<pre>function S'Rounding (X : T) return T</pre>	37
	The function yields the integral value nearest to X , rounding away from zero if X lies exactly halfway between two integers. A zero result has the sign of X when S'Signed_Zeros is True.	38
S'Unbiased_Rounding	S'Unbiased_Rounding denotes a function with the following specification:	39
	<pre>function S'Unbiased_Rounding (X : T) return T</pre>	40
	The function yields the integral value nearest to X , rounding toward the even integer if X lies exactly halfway between two integers. A zero result has the sign of X when S'Signed_Zeros is True.	41
S'Machine_Rounding	S'Machine_Rounding denotes a function with the following specification:	41.1/2
	<pre>function S'Machine_Rounding (X : T) return T</pre>	41.2/2
	The function yields the integral value nearest to X . If X lies exactly halfway between two integers, one of those integers is returned, but which of them is returned is unspecified. A zero result has the sign of X when S'Signed_Zeros is True. This function provides access to the rounding behavior which is most efficient on the target processor.	41.3/2
S'Truncation	S'Truncation denotes a function with the following specification:	42
	<pre>function S'Truncation (X : T) return T</pre>	43
	The function yields the value $\lceil X \rceil$ when X is negative, and $\lfloor X \rfloor$ otherwise. A zero result has the sign of X when S'Signed_Zeros is True.	44
S'Remainder	S'Remainder denotes a function with the following specification:	45
	<pre>function S'Remainder (X, Y : T) return T</pre>	46
	For nonzero Y , let v be the value $X - n \cdot Y$, where n is the integer nearest to the exact value of X/Y ; if $ n - X/Y = 1/2$, then n is chosen to be even. If v is a machine number of the type T , the function yields v ; otherwise, it yields zero. Constraint_Error is raised if Y is zero. A zero result has the sign of X when S'Signed_Zeros is True.	47
S'Adjacent	S'Adjacent denotes a function with the following specification:	48
	<pre>function S'Adjacent (X, Towards : T) return T</pre>	49
	If $Towards = X$, the function yields X ; otherwise, it yields the machine number of the type T adjacent to X in the direction of $Towards$, if that machine number exists. If the result would be outside the base range of S, Constraint_Error is raised. When T'Signed_Zeros is True, a zero result has the sign of X . When $Towards$ is zero, its sign has no bearing on the result.	50
S'Copy_Sign	S'Copy_Sign denotes a function with the following specification:	51

52 **function** S'Copy_Sign (*Value*, *Sign* : *T*)
 return *T*

53 If the value of *Value* is nonzero, the function yields a result whose magnitude is that of *Value* and whose sign is that of *Sign*; otherwise, it yields the value zero. *Constraint_Error* is optionally raised if the result is outside the base range of S. A zero result has the sign of *Sign* when S'Signed_Zeros is True.

54 S'Leading_Part
 S'Leading_Part denotes a function with the following specification:

55 **function** S'Leading_Part (*X* : *T*;
 Radix_Digits : *universal_integer*)
 return *T*

56 Let *v* be the value $T\text{Machine_Radix}^{k-\text{Radix_Digits}}$, where *k* is the normalized exponent of *X*. The function yields the value

- 57 • $\lfloor X/v \rfloor \cdot v$, when *X* is nonnegative and *Radix_Digits* is positive;
- 58 • $\lceil X/v \rceil \cdot v$, when *X* is negative and *Radix_Digits* is positive.

59 *Constraint_Error* is raised when *Radix_Digits* is zero or negative. A zero result, which can only occur when *X* is zero, has the sign of *X*.

60 S'Machine S'Machine denotes a function with the following specification:

61 **function** S'Machine (*X* : *T*)
 return *T*

62 If *X* is a machine number of the type *T*, the function yields *X*; otherwise, it yields the value obtained by rounding or truncating *X* to either one of the adjacent machine numbers of the type *T*. *Constraint_Error* is raised if rounding or truncating *X* to the precision of the machine numbers results in a value outside the base range of S. A zero result has the sign of *X* when S'Signed_Zeros is True.

63 The following *model-oriented attributes* are defined for any subtype S of a floating point type *T*.

64 S'Model_Mantissa

If the Numerics Annex is not supported, this attribute yields an implementation defined value that is greater than or equal to $\lceil d \cdot \log(10) / \log(T\text{Machine_Radix}) \rceil + 1$, where *d* is the requested decimal precision of *T*, and less than or equal to the value of *TMachine_Mantissa*. See G.2.2 for further requirements that apply to implementations supporting the Numerics Annex. The value of this attribute is of the type *universal_integer*.

65 S'Model_Emin

If the Numerics Annex is not supported, this attribute yields an implementation defined value that is greater than or equal to the value of *TMachine_Emin*. See G.2.2 for further requirements that apply to implementations supporting the Numerics Annex. The value of this attribute is of the type *universal_integer*.

66 S'Model_Epsilon

Yields the value $T\text{Machine_Radix}^{1 - T\text{Model_Mantissa}}$. The value of this attribute is of the type *universal_real*.

67 S'Model_Small

Yields the value $T\text{Machine_Radix}^{T\text{Model_Emin} - 1}$. The value of this attribute is of the type *universal_real*.

68 S'Model S'Model denotes a function with the following specification:

69 **function** S'Model (*X* : *T*)
 return *T*

If the Numerics Annex is not supported, the meaning of this attribute is implementation defined; see G.2.2 for the definition that applies to implementations supporting the Numerics Annex.

S'Safe_First

Yields the lower bound of the safe range (see 3.5.7) of the type *T*. If the Numerics Annex is not supported, the value of this attribute is implementation defined; see G.2.2 for the definition that applies to implementations supporting the Numerics Annex. The value of this attribute is of the type *universal_real*.

S'Safe_Last

Yields the upper bound of the safe range (see 3.5.7) of the type *T*. If the Numerics Annex is not supported, the value of this attribute is implementation defined; see G.2.2 for the definition that applies to implementations supporting the Numerics Annex. The value of this attribute is of the type *universal_real*.

A.5.4 Attributes of Fixed Point Types*Static Semantics*

The following *representation-oriented* attributes are defined for every subtype *S* of a fixed point type *T*.

S'Machine_Radix

Yields the radix of the hardware representation of the type *T*. The value of this attribute is of the type *universal_integer*.

S'Machine_Rounds

Yields the value True if rounding is performed on inexact results of every predefined operation that yields a result of the type *T*; yields the value False otherwise. The value of this attribute is of the predefined type Boolean.

S'Machine_Overflows

Yields the value True if overflow and divide-by-zero are detected and reported by raising Constraint_Error for every predefined operation that yields a result of the type *T*; yields the value False otherwise. The value of this attribute is of the predefined type Boolean.

A.6 Input-Output

Input-output is provided through language-defined packages, each of which is a child of the root package Ada. The generic packages Sequential_IO and Direct_IO define input-output operations applicable to files containing elements of a given type. The generic package Storage_IO supports reading from and writing to an in-memory buffer. Additional operations for text input-output are supplied in the packages Text_IO, Wide_Text_IO, and Wide_Wide_Text_IO. Heterogeneous input-output is provided through the child packages Streams.Stream_IO and Text_IO.Text_Streams (see also 13.13). The package IO_Exceptions defines the exceptions needed by the predefined input-output packages.

A.7 External Files and File Objects*Static Semantics*

Values input from the external environment of the program, or output to the external environment, are considered to occupy *external files*. An external file can be anything external to the program that can produce a value to be read or receive a value to be written. An external file is identified by a string (the *name*). A second string (the *form*) gives further system-dependent characteristics that may be associated

with the file, such as the physical organization or access rights. The conventions governing the interpretation of such strings shall be documented.

- 2 Input and output operations are expressed as operations on objects of some *file type*, rather than directly in terms of the external files. In the remainder of this section, the term *file* is always used to refer to a file object; the term *external file* is used otherwise.
 - 3 Input-output for sequential files of values of a single element type is defined by means of the generic package `Sequential_IO`. In order to define sequential input-output for a given element type, an instantiation of this generic unit, with the given type as actual parameter, has to be declared. The resulting package contains the declaration of a file type (called `File_Type`) for files of such elements, as well as the operations applicable to these files, such as the `Open`, `Read`, and `Write` procedures.
 - 4/2 Input-output for direct access files is likewise defined by a generic package called `Direct_IO`. Input-output in human-readable form is defined by the (nongeneric) packages `Text_IO` for Character and String data, `Wide_Text_IO` for `Wide_Character` and `Wide_String` data, and `Wide_Wide_Text_IO` for `Wide_Wide_Character` and `Wide_Wide_String` data. Input-output for files containing streams of elements representing values of possibly different types is defined by means of the (nongeneric) package `Streams.Stream_IO`.
 - 5 Before input or output operations can be performed on a file, the file first has to be associated with an external file. While such an association is in effect, the file is said to be *open*, and otherwise the file is said to be *closed*.
 - 6 The language does not define what happens to external files after the completion of the main program and all the library tasks (in particular, if corresponding files have not been closed). The effect of input-output for access types is unspecified.
 - 7 An open file has a *current mode*, which is a value of one of the following enumeration types:


```
8   type File_Mode is (In_File, Inout_File, Out_File);  -- for Direct_IO
9       These values correspond respectively to the cases where only reading, both reading and writing,
          or only writing are to be performed.
10/2  type File_Mode is (In_File, Out_File, Append_File);
        -- for Sequential_IO, Text_IO, Wide_Text_IO, Wide_Wide_Text_IO, and Stream_IO
11       These values correspond respectively to the cases where only reading, only writing, or only
          appending are to be performed.
12   The mode of a file can be changed.
```
 - 13/2 Several file management operations are common to `Sequential_IO`, `Direct_IO`, `Text_IO`, `Wide_Text_IO`, and `Wide_Wide_Text_IO`. These operations are described in subclause A.8.2 for sequential and direct files. Any additional effects concerning text input-output are described in subclause A.10.2.
 - 14 The exceptions that can be propagated by the execution of an input-output subprogram are defined in the package `IO_Exceptions`; the situations in which they can be propagated are described following the description of the subprogram (and in clause A.13). The exceptions `Storage_Error` and `Program_Error` may be propagated. (`Program_Error` can only be propagated due to errors made by the caller of the subprogram.) Finally, exceptions can be propagated in certain implementation-defined situations.
- NOTES
- 15/2 18 Each instantiation of the generic packages `Sequential_IO` and `Direct_IO` declares a different type `File_Type`. In the case of `Text_IO`, `Wide_Text_IO`, `Wide_Wide_Text_IO`, and `Streams.Stream_IO`, the corresponding type `File_Type` is unique.

19 A bidirectional device can often be modeled as two sequential files associated with the device, one of mode In_File, and one of mode Out_File. An implementation may restrict the number of files that may be associated with a given external file.

16

A.8 Sequential and Direct Files

Static Semantics

Two kinds of access to external files are defined in this subclause: *sequential access* and *direct access*. The corresponding file types and the associated operations are provided by the generic packages Sequential_IO and Direct_IO. A file object to be used for sequential access is called a *sequential file*, and one to be used for direct access is called a *direct file*. Access to *stream files* is described in A.12.1.

1/2

For sequential access, the file is viewed as a sequence of values that are transferred in the order of their appearance (as produced by the program or by the external environment). When the file is opened with mode In_File or Out_File, transfer starts respectively from or to the beginning of the file. When the file is opened with mode Append_File, transfer to the file starts after the last element of the file.

2

For direct access, the file is viewed as a set of elements occupying consecutive positions in linear order; a value can be transferred to or from an element of the file at any selected position. The position of an element is specified by its *index*, which is a number, greater than zero, of the implementation-defined integer type Count. The first element, if any, has index one; the index of the last element, if any, is called the *current size*; the current size is zero if there are no elements. The current size is a property of the external file.

3

An open direct file has a *current index*, which is the index that will be used by the next read or write operation. When a direct file is opened, the current index is set to one. The current index of a direct file is a property of a file object, not of an external file.

4

A.8.1 The Generic Package Sequential_IO

Static Semantics

The generic library package Sequential_IO has the following declaration:

1

```

with Ada.IO_Exceptions;
generic
  type Element_Type(<>) is private;
package Ada.Sequential_IO is
  type File_Type is limited private;
  type File_Mode is (In_File, Out_File, Append_File);
  -- File management
  procedure Create(File : in out File_Type;
                  Mode : in File_Mode := Out_File;
                  Name : in String := "";
                  Form : in String := "");
  procedure Open   (File : in out File_Type;
                  Mode : in File_Mode;
                  Name : in String;
                  Form : in String := "");
  procedure Close (File : in out File_Type);
  procedure Delete(File : in out File_Type);
  procedure Reset (File : in out File_Type; Mode : in File_Mode);
  procedure Reset (File : in out File_Type);

```

2

3

4

5

6

7

8

```

9      function Mode    (File : in File_Type) return File_Mode;
10     function Name    (File : in File_Type) return String;
11     function Form    (File : in File_Type) return String;
12     function Is_Open (File : in File_Type) return Boolean;
13     -- Input and output operations
14
15     procedure Read   (File : in File_Type; Item : out Element_Type);
16     procedure Write  (File : in File_Type; Item : in Element_Type);
17     function End_Of_File (File : in File_Type) return Boolean;
18     -- Exceptions
19
20     Status_Error : exception renames IO_Exceptions.Status_Error;
21     Mode_Error   : exception renames IO_Exceptions.Mode_Error;
22     Name_Error   : exception renames IO_Exceptions.Name_Error;
23     Use_Error    : exception renames IO_Exceptions.Use_Error;
24     Device_Error : exception renames IO_Exceptions.Device_Error;
25     End_Error    : exception renames IO_Exceptions.End_Error;
26     Data_Error   : exception renames IO_Exceptions.Data_Error;
27
28 private
29     . . . . . not specified by the language
30 end Ada.Sequential_IO;

```

17/2 The type `File_Type` needs finalization (see 7.6) in every instantiation of `Sequential_IO`.

A.8.2 File Management

Static Semantics

1 The procedures and functions described in this subclause provide for the control of external files; their declarations are repeated in each of the packages for sequential, direct, text, and stream input-output. For text input-output, the procedures Create, Open, and Reset have additional effects described in subclause A.10.2.

```

2      procedure Create (File : in out File_Type;
3                         Mode : in File_Mode := default_mode;
4                         Name : in String := "";
5                         Form : in String := "");

```

3/2 Establishes a new external file, with the given name and form, and associates this external file with the given file. The given file is left open. The current mode of the given file is set to the given access mode. The default access mode is the mode `Out_File` for sequential, stream, and text input-output; it is the mode `Inout_File` for direct input-output. For direct access, the size of the created file is implementation defined.

4 A null string for `Name` specifies an external file that is not accessible after the completion of the main program (a temporary file). A null string for `Form` specifies the use of the default options of the implementation for the external file.

5 The exception `Status_Error` is propagated if the given file is already open. The exception `Name_Error` is propagated if the string given as `Name` does not allow the identification of an external file. The exception `Use_Error` is propagated if, for the specified mode, the external environment does not support creation of an external file with the given name (in the absence of `Name_Error`) and form.

```
procedure Open(File : in out File_Type;
               Mode : in File_Mode;
               Name : in String;
               Form : in String := "");
```

Associates the given file with an existing external file having the given name and form, and sets the current mode of the given file to the given mode. The given file is left open.

The exception Status_Error is propagated if the given file is already open. The exception Name_Error is propagated if the string given as Name does not allow the identification of an external file; in particular, this exception is propagated if no external file with the given name exists. The exception Use_Error is propagated if, for the specified mode, the external environment does not support opening for an external file with the given name (in the absence of Name_Error) and form.

```
procedure Close(File : in out File_Type);
```

Severs the association between the given file and its associated external file. The given file is left closed. In addition, for sequential files, if the file being closed has mode Out_File or Append_File, then the last element written since the most recent open or reset is the last element that can be read from the file. If no elements have been written and the file mode is Out_File, then the closed file is empty. If no elements have been written and the file mode is Append_File, then the closed file is unchanged.

The exception Status_Error is propagated if the given file is not open.

```
procedure Delete(File : in out File_Type);
```

Deletes the external file associated with the given file. The given file is closed, and the external file ceases to exist.

The exception Status_Error is propagated if the given file is not open. The exception Use_Error is propagated if deletion of the external file is not supported by the external environment.

```
procedure Reset(File : in out File_Type; Mode : in File_Mode);
procedure Reset(File : in out File_Type);
```

Resets the given file so that reading from its elements can be restarted from the beginning of the external file (for modes In_File and Inout_File), and so that writing to its elements can be restarted at the beginning of the external file (for modes Out_File and Inout_File) or after the last element of the external file (for mode Append_File). In particular, for direct access this means that the current index is set to one. If a Mode parameter is supplied, the current mode of the given file is set to the given mode. In addition, for sequential files, if the given file has mode Out_File or Append_File when Reset is called, the last element written since the most recent open or reset is the last element that can be read from the external file. If no elements have been written and the file mode is Out_File, the reset file is empty. If no elements have been written and the file mode is Append_File, then the reset file is unchanged.

The exception Status_Error is propagated if the file is not open. The exception Use_Error is propagated if the external environment does not support resetting for the external file and, also, if the external environment does not support resetting to the specified mode for the external file.

```
function Mode(File : in File_Type) return File_Mode;
```

Returns the current mode of the given file.

The exception Status_Error is propagated if the file is not open.

```

21   function Name(File : in File_Type) return String;
22/2   Returns a string which uniquely identifies the external file currently associated with the given
      file (and may thus be used in an Open operation).
23   The exception Status_Error is propagated if the given file is not open. The exception Use_Error
      is propagated if the associated external file is a temporary file that cannot be opened by any
      name.
24   function Form(File : in File_Type) return String;
25   Returns the form string for the external file currently associated with the given file. If an
      external environment allows alternative specifications of the form (for example, abbreviations
      using default options), the string returned by the function should correspond to a full
      specification (that is, it should indicate explicitly all options selected, including default options).
26   The exception Status_Error is propagated if the given file is not open.
27   function Is_Open(File : in File_Type) return Boolean;
28   Returns True if the file is open (that is, if it is associated with an external file), otherwise returns
      False.

```

Implementation Permissions

29 An implementation may propagate Name_Error or Use_Error if an attempt is made to use an I/O feature that cannot be supported by the implementation due to limitations in the external environment. Any such restriction should be documented.

A.8.3 Sequential Input-Output Operations

Static Semantics

1 The operations available for sequential input and output are described in this subclause. The exception Status_Error is propagated if any of these operations is attempted for a file that is not open.

```

2   procedure Read(File : in File_Type; Item : out Element_Type);
3       Operates on a file of mode In_File. Reads an element from the given file, and returns the value
          of this element in the Item parameter.
4       The exception Mode_Error is propagated if the mode is not In_File. The exception End_Error is
          propagated if no more elements can be read from the given file. The exception Data_Error can
          be propagated if the element read cannot be interpreted as a value of the subtype Element_Type
          (see A.13, “Exceptions in Input-Output”).
5   procedure Write(File : in File_Type; Item : in Element_Type);
6       Operates on a file of mode Out_File or Append_File. Writes the value of Item to the given file.
7       The exception Mode_Error is propagated if the mode is not Out_File or Append_File. The
          exception Use_Error is propagated if the capacity of the external file is exceeded.
8   function End_Of_File(File : in File_Type) return Boolean;
9       Operates on a file of mode In_File. Returns True if no more elements can be read from the given
          file; otherwise returns False.
10      The exception Mode_Error is propagated if the mode is not In_File.

```

A.8.4 The Generic Package Direct_IO

Static Semantics

The generic library package Direct_IO has the following declaration:

```

1      with Ada.IO_Exceptions;
2      generic
3          type Element_Type is private;
4          package Ada.Direct_IO is
5              type File_Type is limited private;
6              type File_Mode is (In_File, Inout_File, Out_File);
7              type Count    is range 0 .. implementation-defined;
8              subtype Positive_Count is Count range 1 .. Count'Last;
9              -- File management
10             procedure Create(File : in out File_Type;
11                             Mode : in File_Mode := Inout_File;
12                             Name : in String := "";
13                             Form : in String := "");
14             procedure Open   (File : in out File_Type;
15                             Mode : in File_Mode;
16                             Name : in String;
17                             Form : in String := "");
18             procedure Close  (File : in out File_Type);
19             procedure Delete (File : in out File_Type);
20             procedure Reset  (File : in out File_Type; Mode : in File_Mode);
21             procedure Reset  (File : in out File_Type);
22             function Mode    (File : in File_Type) return File_Mode;
23             function Name    (File : in File_Type) return String;
24             function Form    (File : in File_Type) return String;
25             function Is_Open (File : in File_Type) return Boolean;
26             -- Input and output operations
27             procedure Read  (File : in File_Type; Item : out Element_Type;
28                             From : in Positive_Count);
29             procedure Read  (File : in File_Type; Item : out Element_Type);
30             procedure Write (File : in File_Type; Item : in Element_Type;
31                             To   : in Positive_Count);
32             procedure Write (File : in File_Type; Item : in Element_Type);
33             procedure Set_Index (File : in File_Type; To : in Positive_Count);
34             function Index (File : in File_Type) return Positive_Count;
35             function Size   (File : in File_Type) return Count;
36             function End_Of_File (File : in File_Type) return Boolean;
37             -- Exceptions
38             Status_Error : exception renames IO_Exceptions.Status_Error;
39             Mode_Error   : exception renames IO_Exceptions.Mode_Error;
40             Name_Error   : exception renames IO_Exceptions.Name_Error;
41             Use_Error    : exception renames IO_Exceptions.Use_Error;
42             Device_Error : exception renames IO_Exceptions.Device_Error;
43             End_Error    : exception renames IO_Exceptions.End_Error;
44             Data_Error   : exception renames IO_Exceptions.Data_Error;
45
46             private
47                 ... -- not specified by the language
48             end Ada.Direct_IO;

```

The type File_Type needs finalization (see 7.6) in every instantiation of Direct_IO.

A.8.5 Direct Input-Output Operations

Static Semantics

1 The operations available for direct input and output are described in this subclause. The exception Status_Error is propagated if any of these operations is attempted for a file that is not open.

2 **procedure** Read(File : **in** File_Type; Item : **out** Element_Type;
 From : **in** Positive_Count);
 procedure Read(File : **in** File_Type; Item : **out** Element_Type);

3 Operates on a file of mode In_File or Inout_File. In the case of the first form, sets the current index of the given file to the index value given by the parameter From. Then (for both forms) returns, in the parameter Item, the value of the element whose position in the given file is specified by the current index of the file; finally, increases the current index by one.

4 The exception Mode_Error is propagated if the mode of the given file is Out_File. The exception End_Error is propagated if the index to be used exceeds the size of the external file. The exception Data_Error can be propagated if the element read cannot be interpreted as a value of the subtype Element_Type (see A.13).

5 **procedure** Write(File : **in** File_Type; Item : **in** Element_Type;
 To : **in** Positive_Count);
 procedure Write(File : **in** File_Type; Item : **in** Element_Type);

6 Operates on a file of mode Inout_File or Out_File. In the case of the first form, sets the index of the given file to the index value given by the parameter To. Then (for both forms) gives the value of the parameter Item to the element whose position in the given file is specified by the current index of the file; finally, increases the current index by one.

7 The exception Mode_Error is propagated if the mode of the given file is In_File. The exception Use_Error is propagated if the capacity of the external file is exceeded.

8 **procedure** Set_Index(File : **in** File_Type; To : **in** Positive_Count);

9 Operates on a file of any mode. Sets the current index of the given file to the given index value (which may exceed the current size of the file).

10 **function** Index(File : **in** File_Type) **return** Positive_Count;

11 Operates on a file of any mode. Returns the current index of the given file.

12 **function** Size(File : **in** File_Type) **return** Count;

13 Operates on a file of any mode. Returns the current size of the external file that is associated with the given file.

14 **function** End_Of_File(File : **in** File_Type) **return** Boolean;

15 Operates on a file of mode In_File or Inout_File. Returns True if the current index exceeds the size of the external file; otherwise returns False.

16 The exception Mode_Error is propagated if the mode of the given file is Out_File.

NOTES

17 20 Append_File mode is not supported for the generic package Direct_IO.

A.9 The Generic Package Storage_IO

The generic package Storage_IO provides for reading from and writing to an in-memory buffer. This generic package supports the construction of user-defined input-output packages.

Static Semantics

The generic library package Storage_IO has the following declaration:

```

with Ada.IO_Exceptions;
with System.Storage_Elements;
generic
  type Element_Type is private;
package Ada.Storage_IO is
  pragma Preelaborate(Storage_IO);
  Buffer_Size : constant System.Storage_Elements.Storage_Count := 4
    implementation-defined;
  subtype Buffer_Type is 5
    System.Storage_Elements.Storage_Array(1..Buffer_Size);
  -- Input and output operations
  procedure Read (Buffer : in Buffer_Type; Item : out Element_Type); 6
  procedure Write(Buffer : out Buffer_Type; Item : in Element_Type); 7
  -- Exceptions
  Data_Error : exception renames IO_Exceptions.Data_Error; 8
end Ada.Storage_IO; 9

```

In each instance, the constant Buffer_Size has a value that is the size (in storage elements) of the buffer required to represent the content of an object of subtype Element_Type, including any implicit levels of indirection used by the implementation. The Read and Write procedures of Storage_IO correspond to the Read and Write procedures of Direct_IO (see A.8.4), but with the content of the Item parameter being read from or written into the specified Buffer, rather than an external file.

NOTES

21 A buffer used for Storage_IO holds only one element at a time; an external file used for Direct_IO holds a sequence of elements.

A.10 Text Input-Output

Static Semantics

This clause describes the package Text_IO, which provides facilities for input and output in human-readable form. Each file is read or written sequentially, as a sequence of characters grouped into lines, and as a sequence of lines grouped into pages. The specification of the package is given below in subclause A.10.1.

The facilities for file management given above, in subclauses A.8.2 and A.8.3, are available for text input-output. In place of Read and Write, however, there are procedures Get and Put that input values of suitable types from text files, and output values to them. These values are provided to the Put procedures, and returned by the Get procedures, in a parameter Item. Several overloaded procedures of these names exist, for different types of Item. These Get procedures analyze the input sequences of characters based on lexical elements (see Section 2) and return the corresponding values; the Put procedures output the given values as appropriate lexical elements. Procedures Get and Put are also available that input and output individual characters treated as character values rather than as lexical elements. Related to character input

are procedures to look ahead at the next character without reading it, and to read a character “immediately” without waiting for an end-of-line to signal availability.

- 3 In addition to the procedures Get and Put for numeric and enumeration types of Item that operate on text files, analogous procedures are provided that read from and write to a parameter of type String. These procedures perform the same analysis and composition of character sequences as their counterparts which have a file parameter.
- 4 For all Get and Put procedures that operate on text files, and for many other subprograms, there are forms with and without a file parameter. Each such Get procedure operates on an input file, and each such Put procedure operates on an output file. If no file is specified, a default input file or a default output file is used.
- 5 At the beginning of program execution the default input and output files are the so-called standard input file and standard output file. These files are open, have respectively the current modes In_File and Out_File, and are associated with two implementation-defined external files. Procedures are provided to change the current default input file and the current default output file.
- 6 At the beginning of program execution a default file for program-dependent error-related text output is the so-called standard error file. This file is open, has the current mode Out_File, and is associated with an implementation-defined external file. A procedure is provided to change the current default error file.
- 7 From a logical point of view, a text file is a sequence of pages, a page is a sequence of lines, and a line is a sequence of characters; the end of a line is marked by a *line terminator*; the end of a page is marked by the combination of a line terminator immediately followed by a *page terminator*; and the end of a file is marked by the combination of a line terminator immediately followed by a page terminator and then a *file terminator*. Terminators are generated during output; either by calls of procedures provided expressly for that purpose; or implicitly as part of other operations, for example, when a bounded line length, a bounded page length, or both, have been specified for a file.
- 8 The actual nature of terminators is not defined by the language and hence depends on the implementation. Although terminators are recognized or generated by certain of the procedures that follow, they are not necessarily implemented as characters or as sequences of characters. Whether they are characters (and if so which ones) in any particular implementation need not concern a user who neither explicitly outputs nor explicitly inputs control characters. The effect of input (Get) or output (Put) of control characters (other than horizontal tabulation) is not specified by the language.
- 9 The characters of a line are numbered, starting from one; the number of a character is called its *column number*. For a line terminator, a column number is also defined: it is one more than the number of characters in the line. The lines of a page, and the pages of a file, are similarly numbered. The current column number is the column number of the next character or line terminator to be transferred. The current line number is the number of the current line. The current page number is the number of the current page. These numbers are values of the subtype Positive_Count of the type Count (by convention, the value zero of the type Count is used to indicate special conditions).
- 10 **type** Count **is range** 0 .. *implementation-defined*;
 subtype Positive_Count **is** Count **range** 1 .. Count'Last;
- 11 For an output file or an append file, a *maximum line length* can be specified and a *maximum page length* can be specified. If a value to be output cannot fit on the current line, for a specified maximum line length, then a new line is automatically started before the value is output; if, further, this new line cannot fit on the current page, for a specified maximum page length, then a new page is automatically started before the value is output. Functions are provided to determine the maximum line length and the maximum page

length. When a file is opened with mode Out_File or Append_File, both values are zero: by convention, this means that the line lengths and page lengths are unbounded. (Consequently, output consists of a single line if the subprograms for explicit control of line and page structure are not used.) The constant Unbounded is provided for this purpose.

A.10.1 The Package Text_IO

Static Semantics

The library package Text_IO has the following declaration:

```

1   with Ada.IO_Exceptions;
2   package Ada.Text_IO is
3     type File_Type is limited private;
4     type File_Mode is (In_File, Out_File, Append_File);
5     type Count is range 0 .. implementation-defined;
6     subtype Positive_Count is Count range 1 .. Count'Last;
7     Unbounded : constant Count := 0; -- line and page length
8     subtype Field      is Integer range 0 .. implementation-defined;
9     subtype Number_Base is Integer range 2 .. 16;
10    type Type_Set is (Lower_Case, Upper_Case);
11    -- File Management
12    procedure Create  (File : in out File_Type;
13                      Mode : in File_Mode := Out_File;
14                      Name : in String      := "";
15                      Form : in String      := "");
16    procedure Open    (File : in out File_Type;
17                      Mode : in File_Mode;
18                      Name : in String;
19                      Form : in String := "");
20    procedure Close   (File : in out File_Type);
21    procedure Delete  (File : in out File_Type);
22    procedure Reset   (File : in out File_Type; Mode : in File_Mode);
23    procedure Reset   (File : in out File_Type);
24    function Mode     (File : in File_Type) return File_Mode;
25    function Name     (File : in File_Type) return String;
26    function Form     (File : in File_Type) return String;
27    function Is_Open  (File : in File_Type) return Boolean;
28    -- Control of default input and output files
29    procedure Set_Input (File : in File_Type);
30    procedure Set_Output(File : in File_Type);
31    procedure Set_Error (File : in File_Type);
32    function Standard_Input  return File_Type;
33    function Standard_Output return File_Type;
34    function Standard_Error  return File_Type;
35    function Current_Input   return File_Type;
36    function Current_Output  return File_Type;
37    function Current_Error   return File_Type;
38    type File_Access is access constant File_Type;
39    function Standard_Input  return File_Access;
40    function Standard_Output return File_Access;
41    function Standard_Error  return File_Access;
42    function Current_Input   return File_Access;
43    function Current_Output  return File_Access;
44    function Current_Error   return File_Access;

```

```

21/1      -- Buffer control
21/2      procedure Flush (File : in File_Type);
21/3      procedure Flush;
22      -- Specification of line and page lengths
23      procedure Set_Line_Length(File : in File_Type; To : in Count);
24      procedure Set_Line_Length(To   : in Count);
25      procedure Set_Page_Length(File : in File_Type; To : in Count);
26      procedure Set_Page_Length(To   : in Count);
27      function Line_Length(File : in File_Type) return Count;
28      function Line_Length return Count;
29      function Page_Length(File : in File_Type) return Count;
30      function Page_Length return Count;
31      -- Column, Line, and Page Control
32      procedure New_Line   (File     : in File_Type;
33                           Spacing : in Positive_Count := 1);
34      procedure New_Line   (Spacing : in Positive_Count := 1);
35      procedure Skip_Line  (File     : in File_Type;
36                           Spacing : in Positive_Count := 1);
37      procedure Skip_Line  (Spacing : in Positive_Count := 1);
38      function End_Of_Line(File : in File_Type) return Boolean;
39      function End_Of_Line return Boolean;
40      procedure New_Page   (File : in File_Type);
41      procedure New_Page;
42      procedure Skip_Page  (File : in File_Type);
43      procedure Skip_Page;
44      function End_Of_Page(File : in File_Type) return Boolean;
45      function End_Of_Page return Boolean;
46      function End_Of_File(File : in File_Type) return Boolean;
47      function End_Of_File return Boolean;
48      procedure Set_Col    (File : in File_Type; To : in Positive_Count);
49      procedure Set_Col    (To   : in Positive_Count);
50      procedure Set_Line   (File : in File_Type; To : in Positive_Count);
51      procedure Set_Line   (To   : in Positive_Count);
52      function Col        (File : in File_Type) return Positive_Count;
53      function Col        return Positive_Count;
54      function Line       (File : in File_Type) return Positive_Count;
55      function Line       return Positive_Count;
56      function Page       (File : in File_Type) return Positive_Count;
57      function Page       return Positive_Count;
58      -- Character Input-Output
59      procedure Get(File : in File_Type; Item : out Character);
60      procedure Get(Item : out Character);
61      procedure Put(File : in File_Type; Item : in Character);
62      procedure Put(Item : in Character);
63      procedure Look_Ahead (File      : in File_Type;
64                            Item      : out Character;
65                            End_Of_Line : out Boolean);
66      procedure Look_Ahead (Item      : out Character;
67                            End_Of_Line : out Boolean);
68      procedure Get_Immediate(File      : in File_Type;
69                               Item      : out Character);
70      procedure Get_Immediate(Item      : out Character);

```

```

procedure Get_Immediate(File      : in File_Type;           45
                        Item       : out Character;
                        Available  : out Boolean);
procedure Get_Immediate(Item      : out Character;           46
                        Available  : out Boolean);

-- String Input-Output

procedure Get(File : in File_Type; Item : out String);          46
procedure Get(Item : out String);                                47
procedure Put(File : in File_Type; Item : in String);           48
procedure Put(Item : in String);
procedure Get_Line(File : in File_Type; Item : out String;        49
                   Last   : out Natural);
procedure Get_Line(Item : out String; Last : out Natural);
function Get_Line(File : in File_Type) return String;           49.1/2
function Get_Line return String;
procedure Put_Line(File : in File_Type; Item : in String);        50
procedure Put_Line(Item : in String);

-- Generic packages for Input-Output of Integer Types            51
generic
  type Num is range <>;
package Integer_IO is
  Default_Width : Field := Num'Width;
  Default_Base  : Number_Base := 10;
  procedure Get(File : in File_Type;           52
                Item   : out Num;
                Width : in Field := 0);
  procedure Get(Item : out Num;               53
                Width : in Field := 0);
  procedure Put(File : in File_Type;           54
                Item   : in Num;
                Width : in Field := Default_Width;
                Base   : in Number_Base := Default_Base);
  procedure Put(Item : in Num;               55
                Width : in Field := Default_Width;
                Base   : in Number_Base := Default_Base);
  procedure Get(From : in String;             56
                Item   : out Num;
                Last   : out Positive);
  procedure Put(To   : out String;             57
                Item   : in Num;
                Base   : in Number_Base := Default_Base);
end Integer_IO;

generic
  type Num is mod <>;
package Modular_IO is
  Default_Width : Field := Num'Width;
  Default_Base  : Number_Base := 10;
  procedure Get(File : in File_Type;           58
                Item   : out Num;
                Width : in Field := 0);
  procedure Get(Item : out Num;               59
                Width : in Field := 0);

```

```

60      procedure Put(File  : in File_Type;
                      Item   : in Num;
                      Width  : in Field := Default_Width;
                      Base   : in Number_Base := Default_Base);
procedure Put(Item  : in Num;
                      Width  : in Field := Default_Width;
                      Base   : in Number_Base := Default_Base);
procedure Get(From : in String;
                      Item   : out Num;
                      Last   : out Positive);
procedure Put(To   : out String;
                      Item   : in Num;
                      Base   : in Number_Base := Default_Base);

61  end Modular_IO;
-- Generic packages for Input-Output of Real Types
62
63 generic
64     type Num is digits <>;
65 package Float_IO is
66     Default_Fore : Field := 2;
       Default_Aft  : Field := Num'Digits-1;
       Default_Exp   : Field := 3;
procedure Get(File  : in File_Type;
                      Item   : out Num;
                      Width  : in Field := 0);
procedure Get(Item  : out Num;
                      Width  : in Field := 0);
procedure Put(File : in File_Type;
                      Item   : in Num;
                      Fore   : in Field := Default_Fore;
                      Aft    : in Field := Default_Aft;
                      Exp    : in Field := Default_Exp);
procedure Put(Item : in Num;
                      Fore   : in Field := Default_Fore;
                      Aft    : in Field := Default_Aft;
                      Exp    : in Field := Default_Exp);
67 procedure Get(From : in String;
                      Item   : out Num;
                      Last   : out Positive);
procedure Put(To   : out String;
                      Item   : in Num;
                      Aft   : in Field := Default_Aft;
                      Exp   : in Field := Default_Exp);
end Float_IO;
68 generic
69     type Num is delta <>;
70 package Fixed_IO is
71     Default_Fore : Field := Num'Fore;
       Default_Aft  : Field := Num'Aft;
       Default_Exp   : Field := 0;
procedure Get(File  : in File_Type;
                      Item   : out Num;
                      Width  : in Field := 0);
procedure Get(Item  : out Num;
                      Width  : in Field := 0);

```

```

procedure Put(File : in File_Type;
              Item : in Num;
              Fore : in Field := Default_Fore;
              Aft : in Field := Default_Aft;
              Exp : in Field := Default_Exp);
procedure Put(Item : in Num;
              Fore : in Field := Default_Fore;
              Aft : in Field := Default_Aft;
              Exp : in Field := Default_Exp);
procedure Get(From : in String;
              Item : out Num;
              Last : out Positive);
procedure Put(To : out String;
              Item : in Num;
              Aft : in Field := Default_Aft;
              Exp : in Field := Default_Exp);
end Fixed_IO;                                         71

generic
  type Num is delta <> digits <>;
package Decimal_IO is
  Default_Fore : Field := Num'Fore;
  Default_Aft : Field := Num'Aft;
  Default_Exp : Field := 0;
  procedure Get(File : in File_Type;
                Item : out Num;
                Width : in Field := 0);
  procedure Get(Item : out Num;
                Width : in Field := 0);
  procedure Put(File : in File_Type;
                Item : in Num;
                Fore : in Field := Default_Fore;
                Aft : in Field := Default_Aft;
                Exp : in Field := Default_Exp);
  procedure Put(Item : in Num;
                Fore : in Field := Default_Fore;
                Aft : in Field := Default_Aft;
                Exp : in Field := Default_Exp);
  procedure Get(From : in String;
                Item : out Num;
                Last : out Positive);
  procedure Put(To : out String;
                Item : in Num;
                Aft : in Field := Default_Aft;
                Exp : in Field := Default_Exp);
end Decimal_IO;                                         72

-- Generic package for Input-Output of Enumeration Types
73
generic
  type Enum is (<>);
package Enumeration_IO is
  Default_Width : Field := 0;
  Default_Setting : Type_Set := Upper_Case;
  procedure Get(File : in File_Type;
                Item : out Enum);
  procedure Get(Item : out Enum);
  procedure Put(File : in File_Type;
                Item : in Enum;
                Width : in Field := Default_Width;
                Set : in Type_Set := Default_Setting);
  procedure Put(Item : in Enum;
                Width : in Field := Default_Width;
                Set : in Type_Set := Default_Setting); 74
 75
 76
 77
 78
 79
 80
 81
 82

```

```

83      procedure Get(From : in String;
                  Item : out Enum;
                  Last : out Positive);
     procedure Put(To : out String;
                  Item : in Enum;
                  Set : in Type_Set := Default_Setting);
end Enumeration_IO;

84 -- Exceptions

85 Status_Error : exception renames IO_Exceptions.Status_Error;
Mode_Error : exception renames IO_Exceptions.Mode_Error;
Name_Error : exception renames IO_Exceptions.Name_Error;
Use_Error : exception renames IO_Exceptions.Use_Error;
Device_Error : exception renames IO_Exceptions.Device_Error;
End_Error : exception renames IO_Exceptions.End_Error;
Data_Error : exception renames IO_Exceptions.Data_Error;
Layout_Error : exception renames IO_Exceptions.Layout_Error;
private
  . . . -- not specified by the language
end Ada.Text_IO;

```

86/2 The type File_Type needs finalization (see 7.6).

A.10.2 Text File Management

Static Semantics

- 1 The only allowed file modes for text files are the modes In_File, Out_File, and Append_File. The subprograms given in subclause A.8.2 for the control of external files, and the function End_Of_File given in subclause A.8.3 for sequential input-output, are also available for text files. There is also a version of End_Of_File that refers to the current default input file. For text files, the procedures have the following additional effects:
 - 2 • For the procedures Create and Open: After a file with mode Out_File or Append_File is opened, the page length and line length are unbounded (both have the conventional value zero). After a file (of any mode) is opened, the current column, current line, and current page numbers are set to one. If the mode is Append_File, it is implementation defined whether a page terminator will separate preexisting text in the file from the new text to be written.
 - 3 • For the procedure Close: If the file has the current mode Out_File or Append_File, has the effect of calling New_Page, unless the current page is already terminated; then outputs a file terminator.
 - 4 • For the procedure Reset: If the file has the current mode Out_File or Append_File, has the effect of calling New_Page, unless the current page is already terminated; then outputs a file terminator. The current column, line, and page numbers are set to one, and the line and page lengths to Unbounded. If the new mode is Append_File, it is implementation defined whether a page terminator will separate preexisting text in the file from the new text to be written.
- 5 The exception Mode_Error is propagated by the procedure Reset upon an attempt to change the mode of a file that is the current default input file, the current default output file, or the current default error file.

NOTES

- 6 22 An implementation can define the Form parameter of Create and Open to control effects including the following:
 - 7 • the interpretation of line and column numbers for an interactive file, and
 - 8 • the interpretation of text formats in a file created by a foreign program.

A.10.3 Default Input, Output, and Error Files

Static Semantics

The following subprograms provide for the control of the particular default files that are used when a file parameter is omitted from a Get, Put, or other operation of text input-output described below, or when application-dependent error-related text is to be output.

```
procedure Set_Input(File : in File_Type);
```

Operates on a file of mode In_File. Sets the current default input file to File.

The exception Status_Error is propagated if the given file is not open. The exception Mode_Error is propagated if the mode of the given file is not In_File.

```
procedure Set_Output(File : in File_Type);
procedure Set_Error (File : in File_Type);
```

Each operates on a file of mode Out_File or Append_File. Set_Output sets the current default output file to File. Set_Error sets the current default error file to File. The exception Status_Error is propagated if the given file is not open. The exception Mode_Error is propagated if the mode of the given file is not Out_File or Append_File.

```
function Standard_Input return File_Type;
function Standard_Input return File_Access;
```

Returns the standard input file (see A.10), or an access value designating the standard input file, respectively.

```
function Standard_Output return File_Type;
function Standard_Output return File_Access;
```

Returns the standard output file (see A.10) or an access value designating the standard output file, respectively.

```
function Standard_Error return File_Type;
function Standard_Error return File_Access;
```

Returns the standard error file (see A.10), or an access value designating the standard error file, respectively.

The Form strings implicitly associated with the opening of Standard_Input, Standard_Output, and Standard_Error at the start of program execution are implementation defined.

```
function Current_Input return File_Type;
function Current_Input return File_Access;
```

Returns the current default input file, or an access value designating the current default input file, respectively.

```
function Current_Output return File_Type;
function Current_Output return File_Access;
```

Returns the current default output file, or an access value designating the current default output file, respectively.

```

18   function Current_Error return File_Type;
19   function Current_Error return File_Access;
20/1  procedure Flush (File : in File_Type);
21      Returns the current default error file, or an access value designating the current default error file,
           respectively.

22/1  procedure Flush;
23/1  The effect of Flush is the same as the corresponding subprogram in Streams.Stream_IO (see
           A.12.1). If File is not explicitly specified, Current_Output is used.

```

Erroneous Execution

- 22/1 The execution of a program is erroneous if it invokes an operation on a current default input, default output, or default error file, and if the corresponding file object is closed or no longer exists.
- 23/1 *This paragraph was deleted.*

NOTES

- 24 23 The standard input, standard output, and standard error files cannot be opened, closed, reset, or deleted, because the parameter File of the corresponding procedures has the mode **in out**.
- 25 24 The standard input, standard output, and standard error files are different file objects, but not necessarily different external files.

A.10.4 Specification of Line and Page Lengths

Static Semantics

- 1 The subprograms described in this subclause are concerned with the line and page structure of a file of mode Out_File or Append_File. They operate either on the file given as the first parameter, or, in the absence of such a file parameter, on the current default output file. They provide for output of text with a specified maximum line length or page length. In these cases, line and page terminators are output implicitly and automatically when needed. When line and page lengths are unbounded (that is, when they have the conventional value zero), as in the case of a newly opened file, new lines and new pages are only started when explicitly called for.
- 2 In all cases, the exception Status_Error is propagated if the file to be used is not open; the exception Mode_Error is propagated if the mode of the file is not Out_File or Append_File.

```

3   procedure Set_Line_Length(File : in File_Type; To : in Count);
4      Sets the maximum line length of the specified output or append file to the number of characters
           specified by To. The value zero for To specifies an unbounded line length.
5      The exception Use_Error is propagated if the specified line length is inappropriate for the
           associated external file.

6   procedure Set_Page_Length(File : in File_Type; To : in Count);
7      Sets the maximum page length of the specified output or append file to the number of lines
           specified by To. The value zero for To specifies an unbounded page length.
8      The exception Use_Error is propagated if the specified page length is inappropriate for the
           associated external file.

```

```
function Line_Length(File : in File_Type) return Count;
function Line_Length return Count;
```

Returns the maximum line length currently set for the specified output or append file, or zero if the line length is unbounded.

```
function Page_Length(File : in File_Type) return Count;
function Page_Length return Count;
```

Returns the maximum page length currently set for the specified output or append file, or zero if the page length is unbounded.

A.10.5 Operations on Columns, Lines, and Pages

Static Semantics

The subprograms described in this subclause provide for explicit control of line and page structure; they operate either on the file given as the first parameter, or, in the absence of such a file parameter, on the appropriate (input or output) current default file. The exception Status_Error is propagated by any of these subprograms if the file to be used is not open.

```
procedure New_Line(File : in File_Type; Spacing : in Positive_Count := 1);
procedure New_Line(Spacing : in Positive_Count := 1);
```

Operates on a file of mode Out_File or Append_File.

For a Spacing of one: Outputs a line terminator and sets the current column number to one. Then increments the current line number by one, except in the case that the current line number is already greater than or equal to the maximum page length, for a bounded page length; in that case a page terminator is output, the current page number is incremented by one, and the current line number is set to one.

For a Spacing greater than one, the above actions are performed Spacing times.

The exception Mode_Error is propagated if the mode is not Out_File or Append_File.

```
procedure Skip_Line(File : in File_Type; Spacing : in Positive_Count := 1);
procedure Skip_Line(Spacing : in Positive_Count := 1);
```

Operates on a file of mode In_File.

For a Spacing of one: Reads and discards all characters until a line terminator has been read, and then sets the current column number to one. If the line terminator is not immediately followed by a page terminator, the current line number is incremented by one. Otherwise, if the line terminator is immediately followed by a page terminator, then the page terminator is skipped, the current page number is incremented by one, and the current line number is set to one.

For a Spacing greater than one, the above actions are performed Spacing times.

The exception Mode_Error is propagated if the mode is not In_File. The exception End_Error is propagated if an attempt is made to read a file terminator.

```
function End_Of_Line(File : in File_Type) return Boolean;
function End_Of_Line return Boolean;
```

Operates on a file of mode In_File. Returns True if a line terminator or a file terminator is next; otherwise returns False.

The exception Mode_Error is propagated if the mode is not In_File.

```

15  procedure New_Page(File : in File_Type);
procedure New_Page;
16  Operates on a file of mode Out_File or Append_File. Outputs a line terminator if the current line
   is not terminated, or if the current page is empty (that is, if the current column and line numbers
   are both equal to one). Then outputs a page terminator, which terminates the current page. Adds
   one to the current page number and sets the current column and line numbers to one.
17  The exception Mode_Error is propagated if the mode is not Out_File or Append_File.

18  procedure Skip_Page(File : in File_Type);
procedure Skip_Page;
19  Operates on a file of mode In_File. Reads and discards all characters and line terminators until a
   page terminator has been read. Then adds one to the current page number, and sets the current
   column and line numbers to one.
20  The exception Mode_Error is propagated if the mode is not In_File. The exception End_Error is
   propagated if an attempt is made to read a file terminator.

21  function End_Of_Page(File : in File_Type) return Boolean;
function End_Of_Page return Boolean;
22  Operates on a file of mode In_File. Returns True if the combination of a line terminator and a
   page terminator is next, or if a file terminator is next; otherwise returns False.
23  The exception Mode_Error is propagated if the mode is not In_File.

24  function End_Of_File(File : in File_Type) return Boolean;
function End_Of_File return Boolean;
25  Operates on a file of mode In_File. Returns True if a file terminator is next, or if the
   combination of a line, a page, and a file terminator is next; otherwise returns False.
26  The exception Mode_Error is propagated if the mode is not In_File.

27  The following subprograms provide for the control of the current position of reading or writing in a file. In
   all cases, the default file is the current output file.

28  procedure Set_Col(File : in File_Type; To : in Positive_Count);
procedure Set_Col(To : in Positive_Count);
29  If the file mode is Out_File or Append_File:
30    • If the value specified by To is greater than the current column number, outputs spaces,
       adding one to the current column number after each space, until the current column
       number equals the specified value. If the value specified by To is equal to the current
       column number, there is no effect. If the value specified by To is less than the current
       column number, has the effect of calling New_Line (with a spacing of one), then
       outputs (To – 1) spaces, and sets the current column number to the specified value.
31    • The exception Layout_Error is propagated if the value specified by To exceeds
       Line_Length when the line length is bounded (that is, when it does not have the
       conventional value zero).

32  If the file mode is In_File:
33    • Reads (and discards) individual characters, line terminators, and page terminators,
       until the next character to be read has a column number that equals the value specified
       by To; there is no effect if the current column number already equals this value. Each
       transfer of a character or terminator maintains the current column, line, and page

```

numbers in the same way as a Get procedure (see A.10.6). (Short lines will be skipped until a line is reached that has a character at the specified column position.)

- The exception End_Error is propagated if an attempt is made to read a file terminator.

```
procedure Set_Line(File : in File_Type; To : in Positive_Count);  
procedure Set_Line(To : in Positive_Count);
```

If the file mode is Out_File or Append_File:

- If the value specified by To is greater than the current line number, has the effect of repeatedly calling New_Line (with a spacing of one), until the current line number equals the specified value. If the value specified by To is equal to the current line number, there is no effect. If the value specified by To is less than the current line number, has the effect of calling New_Page followed by a call of New_Line with a spacing equal to (To – 1).
- The exception Layout_Error is propagated if the value specified by To exceeds Page_Length when the page length is bounded (that is, when it does not have the conventional value zero).

If the mode is In_File:

- Has the effect of repeatedly calling Skip_Line (with a spacing of one), until the current line number equals the value specified by To; there is no effect if the current line number already equals this value. (Short pages will be skipped until a page is reached that has a line at the specified line position.)
- The exception End_Error is propagated if an attempt is made to read a file terminator.

```
function Col(File : in File_Type) return Positive_Count;  
function Col return Positive_Count;
```

Returns the current column number.

The exception Layout_Error is propagated if this number exceeds Count'Last.

```
function Line(File : in File_Type) return Positive_Count;  
function Line return Positive_Count;
```

Returns the current line number.

The exception Layout_Error is propagated if this number exceeds Count'Last.

```
function Page(File : in File_Type) return Positive_Count;  
function Page return Positive_Count;
```

Returns the current page number.

The exception Layout_Error is propagated if this number exceeds Count'Last.

The column number, line number, or page number are allowed to exceed Count'Last (as a consequence of the input or output of sufficiently many characters, lines, or pages). These events do not cause any exception to be propagated. However, a call of Col, Line, or Page propagates the exception Layout_Error if the corresponding number exceeds Count'Last.

NOTES

25 A page terminator is always skipped whenever the preceding line terminator is skipped. An implementation may represent the combination of these terminators by a single character, provided that it is properly recognized on input.

A.10.6 Get and Put Procedures

Static Semantics

- 1 The procedures Get and Put for items of the type Character, String, numeric types, and enumeration types are described in subsequent subclauses. Features of these procedures that are common to most of these types are described in this subclause. The Get and Put procedures for items of type Character and String deal with individual character values; the Get and Put procedures for numeric and enumeration types treat the items as lexical elements.
- 2 All procedures Get and Put have forms with a file parameter, written first. Where this parameter is omitted, the appropriate (input or output) current default file is understood to be specified. Each procedure Get operates on a file of mode In_File. Each procedure Put operates on a file of mode Out_File or Append_File.
- 3 All procedures Get and Put maintain the current column, line, and page numbers of the specified file: the effect of each of these procedures upon these numbers is the result of the effects of individual transfers of characters and of individual output or skipping of terminators. Each transfer of a character adds one to the current column number. Each output of a line terminator sets the current column number to one and adds one to the current line number. Each output of a page terminator sets the current column and line numbers to one and adds one to the current page number. For input, each skipping of a line terminator sets the current column number to one and adds one to the current line number; each skipping of a page terminator sets the current column and line numbers to one and adds one to the current page number. Similar considerations apply to the procedures Get_Line, Put_Line, and Set_Col.
- 4 Several Get and Put procedures, for numeric and enumeration types, have *format* parameters which specify field lengths; these parameters are of the nonnegative subtype Field of the type Integer.
- 5/2 Input-output of enumeration values uses the syntax of the corresponding lexical elements. Any Get procedure for an enumeration type begins by skipping any leading blanks, or line or page terminators. A *blank* is defined as a space or a horizontal tabulation character. Next, characters are input only so long as the sequence input is an initial sequence of an identifier or of a character literal (in particular, input ceases when a line terminator is encountered). The character or line terminator that causes input to cease remains available for subsequent input.
- 6 For a numeric type, the Get procedures have a format parameter called Width. If the value given for this parameter is zero, the Get procedure proceeds in the same manner as for enumeration types, but using the syntax of numeric literals instead of that of enumeration literals. If a nonzero value is given, then exactly Width characters are input, or the characters up to a line terminator, whichever comes first; any skipped leading blanks are included in the count. The syntax used for numeric literals is an extended syntax that allows a leading sign (but no intervening blanks, or line or page terminators) and that also allows (for real types) an integer literal as well as forms that have digits only before the point or only after the point.
- 7 Any Put procedure, for an item of a numeric or an enumeration type, outputs the value of the item as a numeric literal, identifier, or character literal, as appropriate. This is preceded by leading spaces if required by the format parameters Width or Fore (as described in later subclauses), and then a minus sign for a negative value; for an enumeration type, the spaces follow instead of leading. The format given for a Put procedure is overridden if it is insufficiently wide, by using the minimum needed width.
- 8 Two further cases arise for Put procedures for numeric and enumeration types, if the line length of the specified output file is bounded (that is, if it does not have the conventional value zero). If the number of characters to be output does not exceed the maximum line length, but is such that they cannot fit on the

current line, starting from the current column, then (in effect) New_Line is called (with a spacing of one) before output of the item. Otherwise, if the number of characters exceeds the maximum line length, then the exception Layout_Error is propagated and nothing is output.

The exception Status_Error is propagated by any of the procedures Get, Get_Line, Put, and Put_Line if the file to be used is not open. The exception Mode_Error is propagated by the procedures Get and Get_Line if the mode of the file to be used is not In_File; and by the procedures Put and Put_Line, if the mode is not Out_File or Append_File.

The exception End_Error is propagated by a Get procedure if an attempt is made to skip a file terminator. The exception Data_Error is propagated by a Get procedure if the sequence finally input is not a lexical element corresponding to the type, in particular if no characters were input; for this test, leading blanks are ignored; for an item of a numeric type, when a sign is input, this rule applies to the succeeding numeric literal. The exception Layout_Error is propagated by a Put procedure that outputs to a parameter of type String, if the length of the actual string is insufficient for the output of the item.

Examples

In the examples, here and in subclauses A.10.8 and A.10.9, the string quotes and the lower case letter b are not transferred; they are shown only to reveal the layout and spaces.

N : Integer;	...	
Get(N);		
-- Characters at input	Sequence input	Value of N
-- bb-12535b	-12535	-12535
-- bb12_535e1b	12_535e1	125350
-- bb12_535e;	12_535e	(none) Data_Error raised

Example of overridden width parameter:

```
Put(Item => -23, Width => 2); -- "-23"
```

A.10.7 Input-Output of Characters and Strings

Static Semantics

For an item of type Character the following procedures are provided:

```
procedure Get(File : in File_Type; Item : out Character);
procedure Get(Item : out Character);
```

After skipping any line terminators and any page terminators, reads the next character from the specified input file and returns the value of this character in the out parameter Item.

The exception End_Error is propagated if an attempt is made to skip a file terminator.

```
procedure Put(File : in File_Type; Item : in Character);
procedure Put(Item : in Character);
```

If the line length of the specified output file is bounded (that is, does not have the conventional value zero), and the current column number exceeds it, has the effect of calling New_Line with a spacing of one. Then, or otherwise, outputs the given character to the file.

```

7      procedure Look_Ahead (File       : in File_Type;
                      Item        : out Character;
                      End_Of_Line : out Boolean);
procedure Look_Ahead (Item       : out Character;
                      End_Of_Line : out Boolean);

8/1    Mode_Error is propagated if the mode of the file is not In_File. Sets End_Of_Line to True if at
      end of line, including if at end of page or at end of file; in each of these cases the value of Item is
      not specified. Otherwise End_Of_Line is set to False and Item is set to the next character
      (without consuming it) from the file.

9      procedure Get_Immediate(File : in File_Type;
                           Item  : out Character);
procedure Get_Immediate(Item : out Character);

10   Reads the next character, either control or graphic, from the specified File or the default input
      file. Mode_Error is propagated if the mode of the file is not In_File. End_Error is propagated if
      at the end of the file. The current column, line and page numbers for the file are not affected.

11   procedure Get_Immediate(File       : in File_Type;
                           Item        : out Character;
                           Available  : out Boolean);
procedure Get_Immediate(Item       : out Character;
                           Available  : out Boolean);

12   If a character, either control or graphic, is available from the specified File or the default input
      file, then the character is read; Available is True and Item contains the value of this character. If
      a character is not available, then Available is False and the value of Item is not specified.
      Mode_Error is propagated if the mode of the file is not In_File. End_Error is propagated if at the
      end of the file. The current column, line and page numbers for the file are not affected.

13/2  For an item of type String the following subprograms are provided:

14   procedure Get(File : in File_Type; Item : out String);
procedure Get(Item : out String);

15   Determines the length of the given string and attempts that number of Get operations for
      successive characters of the string (in particular, no operation is performed if the string is null).

16   procedure Put(File : in File_Type; Item : in String);
procedure Put(Item : in String);

17   Determines the length of the given string and attempts that number of Put operations for
      successive characters of the string (in particular, no operation is performed if the string is null).

17.1/2 function Get_Line(File : in File_Type) return String;
function Get_Line return String;

17.2/2 Returns a result string constructed by reading successive characters from the specified input file,
      and assigning them to successive characters of the result string. The result string has a lower
      bound of 1 and an upper bound of the number of characters read. Reading stops when the end of
      the line is met; Skip_Line is then (in effect) called with a spacing of 1.

17.3/2 Constraint_Error is raised if the length of the line exceeds Positive'Last; in this case, the line
      number and page number are unchanged, and the column number is unspecified but no less than
      it was before the call. The exception End_Error is propagated if an attempt is made to skip a file
      terminator.

```

```
procedure Get_Line(File : in File_Type;
                    Item : out String;
                    Last : out Natural);
procedure Get_Line(Item : out String; Last : out Natural);
```

Reads successive characters from the specified input file and assigns them to successive characters of the specified string. Reading stops if the end of the string is met. Reading also stops if the end of the line is met before meeting the end of the string; in this case Skip_Line is (in effect) called with a spacing of 1. The values of characters not assigned are not specified.

If characters are read, returns in Last the index value such that Item(Last) is the last character assigned (the index of the first character assigned is Item'First). If no characters are read, returns in Last an index value that is one less than Item'First. The exception End_Error is propagated if an attempt is made to skip a file terminator.

```
procedure Put_Line(File : in File_Type; Item : in String);
procedure Put_Line(Item : in String);
```

Calls the procedure Put for the given string, and then the procedure New_Line with a spacing of one.

Implementation Advice

The Get_Immediate procedures should be implemented with unbuffered input. For a device such as a keyboard, input should be “available” if a key has already been typed, whereas for a disk file, input should always be available except at end of file. For a file associated with a keyboard-like device, any line-editing features of the underlying operating system should be disabled during the execution of Get_Immediate.

NOTES

26 Get_Immediate can be used to read a single key from the keyboard “immediately”; that is, without waiting for an end of line. In a call of Get_Immediate without the parameter Available, the caller will wait until a character is available.

27 In a literal string parameter of Put, the enclosing string bracket characters are not output. Each doubled string bracket character in the enclosed string is output as a single string bracket character, as a consequence of the rule for string literals (see 2.6).

28 A string read by Get or written by Put can extend over several lines. An implementation is allowed to assume that certain external files do not contain page terminators, in which case Get_Line and Skip_Line can return as soon as a line terminator is read.

A.10.8 Input-Output for Integer Types

Static Semantics

The following procedures are defined in the generic packages Integer_IO and Modular_IO, which have to be instantiated for the appropriate signed integer or modular type respectively (indicated by Num in the specifications).

Values are output as decimal or based literals, without low line characters or exponent, and, for Integer_IO, preceded by a minus sign if negative. The format (which includes any leading spaces and minus sign) can be specified by an optional field width parameter. Values of widths of fields in output formats are of the nonnegative integer subtype Field. Values of bases are of the integer subtype Number_Base.

```
subtype Number_Base is Integer range 2 .. 16;
```

The default field width and base to be used by output procedures are defined by the following variables that are declared in the generic packages Integer_IO and Modular_IO:

5 Default_Width : Field := Num'Width;
 Default_Base : Number_Base := 10;

6 The following procedures are provided:

```
7       procedure Get(File : in File_Type; Item : out Num; Width : in Field := 0);  

        procedure Get(Item : out Num; Width : in Field := 0);
```

8 If the value of the parameter Width is zero, skips any leading blanks, line terminators, or page terminators, then reads a plus sign if present or (for a signed type only) a minus sign if present, then reads the longest possible sequence of characters matching the syntax of a numeric literal without a point. If a nonzero value of Width is supplied, then exactly Width characters are input, or the characters (possibly none) up to a line terminator, whichever comes first; any skipped leading blanks are included in the count.

9 Returns, in the parameter Item, the value of type Num that corresponds to the sequence input.

10 The exception Data_Error is propagated if the sequence of characters read does not form a legal integer literal or if the value obtained is not of the subtype Num (for Integer_IO) or is not in the base range of Num (for Modular_IO).

```
11     procedure Put(File : in File_Type;  

           Item : in Num;  

           Width : in Field := Default_Width;  

           Base : in Number_Base := Default_Base);  
  

  procedure Put(Item : in Num;  

           Width : in Field := Default_Width;  

           Base : in Number_Base := Default_Base);
```

12 Outputs the value of the parameter Item as an integer literal, with no low lines, no exponent, and no leading zeros (but a single zero for the value zero), and a preceding minus sign for a negative value.

13 If the resulting sequence of characters to be output has fewer than Width characters, then leading spaces are first output to make up the difference.

14 Uses the syntax for decimal literal if the parameter Base has the value ten (either explicitly or through Default_Base); otherwise, uses the syntax for based literal, with any letters in upper case.

```
15     procedure Get(From : in String; Item : out Num; Last : out Positive);
```

16 Reads an integer value from the beginning of the given string, following the same rules as the Get procedure that reads an integer value from a file, but treating the end of the string as a file terminator. Returns, in the parameter Item, the value of type Num that corresponds to the sequence input. Returns in Last the index value such that From(Last) is the last character read.

17 The exception Data_Error is propagated if the sequence input does not have the required syntax or if the value obtained is not of the subtype Num.

```
18     procedure Put(To : out String;  

           Item : in Num;  

           Base : in Number_Base := Default_Base);
```

19 Outputs the value of the parameter Item to the given string, following the same rule as for output to a file, using the length of the given string as the value for Width.

20 Integer_Text_IO is a library package that is a nongeneric equivalent to Text_IO.Integer_IO for the predefined type Integer:

```
with Ada.Text_IO;
package Ada.Integer_Text_IO is new Ada.Text_IO.Integer_IO(Integer);
```

For each predefined signed integer type, a nongeneric equivalent to Text_IO.Integer_IO is provided, with names such as Ada.Long_Integer_Text_IO.

Implementation Permissions

The nongeneric equivalent packages may, but need not, be actual instantiations of the generic package for the appropriate predefined type.

NOTES

29 For Modular_IO, execution of Get propagates Data_Error if the sequence of characters read forms an integer literal outside the range 0..Num'Last.

Examples

This paragraph was deleted.

```
package Int_IO is new Integer_IO(Small_Int); use Int_IO;
-- default format used at instantiation,
-- Default_Width = 4, Default_Base = 10
Put(126);                                -- "b126"
Put(-126, 7);                            -- "bbb-126"
Put(126, Width => 13, Base => 2);      -- "bbb2#1111110#"
```

A.10.9 Input-Output for Real Types

Static Semantics

The following procedures are defined in the generic packages Float_IO, Fixed_IO, and Decimal_IO, which have to be instantiated for the appropriate floating point, ordinary fixed point, or decimal fixed point type respectively (indicated by Num in the specifications).

Values are output as decimal literals without low line characters. The format of each value output consists of a Fore field, a decimal point, an Aft field, and (if a nonzero Exp parameter is supplied) the letter E and an Exp field. The two possible formats thus correspond to:

Fore . Aft

and to:

Fore . Aft E Exp

without any spaces between these fields. The Fore field may include leading spaces, and a minus sign for negative values. The Aft field includes only decimal digits (possibly with trailing zeros). The Exp field includes the sign (plus or minus) and the exponent (possibly with leading zeros).

For floating point types, the default lengths of these fields are defined by the following variables that are declared in the generic package Float_IO:

```
Default_Fore : Field := 2;
Default_Aft  : Field := Num'Digits-1;
Default_Exp   : Field := 3;
```

For ordinary or decimal fixed point types, the default lengths of these fields are defined by the following variables that are declared in the generic packages Fixed_IO and Decimal_IO, respectively:

```
Default_Fore : Field := Num'Fore;
Default_Aft  : Field := Num'Aft;
Default_Exp   : Field := 0;
```

11 The following procedures are provided:

```
12 procedure Get(File : in File_Type; Item : out Num; Width : in Field := 0);
procedure Get(Item : out Num; Width : in Field := 0);
```

13 If the value of the parameter Width is zero, skips any leading blanks, line terminators, or page terminators, then reads the longest possible sequence of characters matching the syntax of any of the following (see 2.4):

- 14 • [+−]numeric_literal
- 15 • [+−]numeral.[exponent]
- 16 • [+−].numeral[exponent]
- 17 • [+−]base#based_numeral.#[exponent]
- 18 • [+−]base#.based_numeral#[exponent]

19 If a nonzero value of Width is supplied, then exactly Width characters are input, or the characters (possibly none) up to a line terminator, whichever comes first; any skipped leading blanks are included in the count.

20 Returns in the parameter Item the value of type Num that corresponds to the sequence input, preserving the sign (positive if none has been specified) of a zero value if Num is a floating point type and Num'Signed_Zeros is True.

21 The exception Data_Error is propagated if the sequence input does not have the required syntax or if the value obtained is not of the subtype Num.

```
22 procedure Put(File : in File_Type;
               Item : in Num;
               Fore : in Field := Default_Fore;
               Aft : in Field := Default_Aft;
               Exp : in Field := Default_Exp);

procedure Put(Item : in Num;
             Fore : in Field := Default_Fore;
             Aft : in Field := Default_Aft;
             Exp : in Field := Default_Exp);
```

23 Outputs the value of the parameter Item as a decimal literal with the format defined by Fore, Aft and Exp. If the value is negative, or if Num is a floating point type where Num'Signed_Zeros is True and the value is a negatively signed zero, then a minus sign is included in the integer part. If Exp has the value zero, then the integer part to be output has as many digits as are needed to represent the integer part of the value of Item, overriding Fore if necessary, or consists of the digit zero if the value of Item has no integer part.

24 If Exp has a value greater than zero, then the integer part to be output has a single digit, which is nonzero except for the value 0.0 of Item.

25 In both cases, however, if the integer part to be output has fewer than Fore characters, including any minus sign, then leading spaces are first output to make up the difference. The number of digits of the fractional part is given by Aft, or is one if Aft equals zero. The value is rounded; a value of exactly one half in the last place is rounded away from zero.

26 If Exp has the value zero, there is no exponent part. If Exp has a value greater than zero, then the exponent part to be output has as many digits as are needed to represent the exponent part of the value of Item (for which a single digit integer part is used), and includes an initial sign (plus or minus). If the exponent part to be output has fewer than Exp characters, including the sign, then

leading zeros precede the digits, to make up the difference. For the value 0.0 of Item, the exponent has the value zero.

```
procedure Get(From : in String; Item : out Num; Last : out Positive);
```

27

Reads a real value from the beginning of the given string, following the same rule as the Get procedure that reads a real value from a file, but treating the end of the string as a file terminator. Returns, in the parameter Item, the value of type Num that corresponds to the sequence input. Returns in Last the index value such that From(Last) is the last character read.

28

The exception Data_Error is propagated if the sequence input does not have the required syntax, or if the value obtained is not of the subtype Num.

29

```
procedure Put(To : out String;
```

30

```
    Item : in Num;
    Aft : in Field := Default_Aft;
    Exp : in Field := Default_Exp);
```

Outputs the value of the parameter Item to the given string, following the same rule as for output to a file, using a value for Fore such that the sequence of characters output exactly fills the string, including any leading spaces.

31

Float_Text_IO is a library package that is a nongeneric equivalent to Text_IO.Float_IO for the predefined type Float:

32

```
with Ada.Text_IO;
package Ada.Float_Text_IO is new Ada.Text_IO.Float_IO(Float);
```

33

For each predefined floating point type, a nongeneric equivalent to Text_IO.Float_IO is provided, with names such as Ada.Long_Float_Text_IO.

34

Implementation Permissions

An implementation may extend Get and Put for floating point types to support special values such as infinities and NaNs.

35

The implementation of Put need not produce an output value with greater accuracy than is supported for the base subtype. The additional accuracy, if any, of the value produced by Put when the number of requested digits in the integer and fractional parts exceeds the required accuracy is implementation defined.

36

The nongeneric equivalent packages may, but need not, be actual instantiations of the generic package for the appropriate predefined type.

37

NOTES

30 For an item with a positive value, if output to a string exactly fills the string without leading spaces, then output of the corresponding negative value will propagate Layout_Error.

38

31 The rules for the Value attribute (see 3.5) and the rules for Get are based on the same set of formats.

39

Examples

This paragraph was deleted.

40/1

```
package Real_IO is new Float_IO(Real); use Real_IO;
-- default format used at instantiation, Default_Exp = 3
```

41

```
X : Real := -123.4567; -- digits 8 (see 3.5.7)
```

42

```
Put(X); -- default format
```

43

```
Put(X, Fore => 5, Aft => 3, Exp => 2); -- "bbb-1.235E+2"
```

43

```
Put(X, 5, 3, 0); -- "b-123.457"
```

A.10.10 Input-Output for Enumeration Types

Static Semantics

- 1 The following procedures are defined in the generic package Enumeration_IO, which has to be instantiated for the appropriate enumeration type (indicated by `Enum` in the specification).
- 2 Values are output using either upper or lower case letters for identifiers. This is specified by the parameter `Set`, which is of the enumeration type `Type_Set`.

3 `type Type_Set is (Lower_Case, Upper_Case);`

- 4 The format (which includes any trailing spaces) can be specified by an optional field width parameter. The default field width and letter case are defined by the following variables that are declared in the generic package Enumeration_IO:

5 `Default_Width : Field := 0;`
 `Default_Setting : Type_Set := Upper_Case;`

- 6 The following procedures are provided:

7 `procedure Get(File : in File_Type; Item : out Enum);`
 `procedure Get(Item : out Enum);`

- 8 After skipping any leading blanks, line terminators, or page terminators, reads an identifier according to the syntax of this lexical element (lower and upper case being considered equivalent), or a character literal according to the syntax of this lexical element (including the apostrophes). Returns, in the parameter `Item`, the value of type `Enum` that corresponds to the sequence input.

- 9 The exception `Data_Error` is propagated if the sequence input does not have the required syntax, or if the identifier or character literal does not correspond to a value of the subtype `Enum`.

10 `procedure Put(File : in File_Type;`
 `Item : in Enum;`
 `Width : in Field := Default_Width;`
 `Set : in Type_Set := Default_Setting);`

`procedure Put(Item : in Enum;`
 `Width : in Field := Default_Width;`
 `Set : in Type_Set := Default_Setting);`

- 11 Outputs the value of the parameter `Item` as an enumeration literal (either an identifier or a character literal). The optional parameter `Set` indicates whether lower case or upper case is used for identifiers; it has no effect for character literals. If the sequence of characters produced has fewer than `Width` characters, then trailing spaces are finally output to make up the difference. If `Enum` is a character type, the sequence of characters produced is as for `Enum'Image(Item)`, as modified by the `Width` and `Set` parameters.

12 `procedure Get(From : in String; Item : out Enum; Last : out Positive);`

- 13 Reads an enumeration value from the beginning of the given string, following the same rule as the `Get` procedure that reads an enumeration value from a file, but treating the end of the string as a file terminator. Returns, in the parameter `Item`, the value of type `Enum` that corresponds to the sequence input. Returns in `Last` the index value such that `From(Last)` is the last character read.

- 14 The exception `Data_Error` is propagated if the sequence input does not have the required syntax, or if the identifier or character literal does not correspond to a value of the subtype `Enum`.

```
procedure Put(To    : out String;
              Item   : in Enum;
              Set    : in Type_Set := Default_Setting);
```

Outputs the value of the parameter Item to the given string, following the same rule as for output to a file, using the length of the given string as the value for Width.

Although the specification of the generic package Enumeration_IO would allow instantiation for an integer type, this is not the intended purpose of this generic package, and the effect of such instantiations is not defined by the language.

NOTES

32 There is a difference between Put defined for characters, and for enumeration values. Thus

Ada.Text_IO.Put('A'); -- outputs the character A

package Char_IO **is new** Ada.Text_IO.Enumeration_IO(Character);
Char_IO.Put('A'); -- outputs the character 'A', between apostrophes

33 The type Boolean is an enumeration type, hence Enumeration_IO can be instantiated for this type.

A.10.11 Input-Output for Bounded Strings

The package Text_IO.Bounded_IO provides input-output in human-readable form for Bounded_Strings.

Static Semantics

The generic library package Text_IO.Bounded_IO has the following declaration:

```
with Ada.Strings.Bounded;
generic
  with package Bounded is
    new Ada.Strings.Bounded.Generic_Bounded_Length(<>);
package Ada.Text_IO.Bounded_IO is
  procedure Put
    (File : in File_Type;
     Item : in Bounded.Bounded_String);
  procedure Put
    (Item : in Bounded.Bounded_String);
  procedure Put_Line
    (File : in File_Type;
     Item : in Bounded.Bounded_String);
  procedure Put_Line
    (Item : in Bounded.Bounded_String);
  function Get_Line
    (File : in File_Type)
    return Bounded.Bounded_String;
  function Get_Line
    return Bounded.Bounded_String;
  procedure Get_Line
    (File : in File_Type; Item : out Bounded.Bounded_String);
  procedure Get_Line
    (Item : out Bounded.Bounded_String);
end Ada.Text_IO.Bounded_IO;
```

For an item of type Bounded_String, the following subprograms are provided:

```
procedure Put
  (File : in File_Type;
   Item : in Bounded.Bounded_String);
```

Equivalent to Text_IO.Put(File, Bounded.To_String(Item));

```

16/2   procedure Put
      (Item : in Bounded.Bounded_String);
17/2     Equivalent to Text_IO.Put (Bounded.To_String(Item));
18/2   procedure Put_Line
      (File : in File_Type;
       Item : in Bounded.Bounded_String);
19/2     Equivalent to Text_IO.Put_Line (File, Bounded.To_String(Item));
20/2   procedure Put_Line
      (Item : in Bounded.Bounded_String);
21/2     Equivalent to Text_IO.Put_Line (Bounded.To_String(Item));
22/2   function Get_Line
      (File : in File_Type)
      return Bounded.Bounded_String;
23/2     Returns Bounded.To_Bounded_String(Text_IO.Get_Line(File));
24/2   function Get_Line
      return Bounded.Bounded_String;
25/2     Returns Bounded.To_Bounded_String(Text_IO.Get_Line);
26/2   procedure Get_Line
      (File : in File_Type; Item : out Bounded.Bounded_String);
27/2     Equivalent to Item := Get_Line (File);
28/2   procedure Get_Line
      (Item : out Bounded.Bounded_String);
29/2     Equivalent to Item := Get_Line;

```

A.10.12 Input-Output for Unbounded Strings

1/2 The package Text_IO.Unbounded_IO provides input-output in human-readable form for Unbounded_Strings.

Static Semantics

```

2/2 The library package Text_IO.Unbounded_IO has the following declaration:
3/2   with Ada.Strings.Unbounded;
4/2   package Ada.Text_IO.Unbounded_IO is
5/2     procedure Put
        (File : in File_Type;
         Item : in Strings.Unbounded.Unbounded_String);
5/2     procedure Put
        (Item : in Strings.Unbounded.Unbounded_String);
6/2     procedure Put_Line
        (File : in File_Type;
         Item : in Strings.Unbounded.Unbounded_String);
7/2     procedure Put_Line
        (Item : in Strings.Unbounded.Unbounded_String);
8/2     function Get_Line
        (File : in File_Type)
        return Strings.Unbounded.Unbounded_String;
9/2     function Get_Line
        return Strings.Unbounded.Unbounded_String;

```

```

procedure Get_Line          10/2
  (File : in File_Type; Item : out Strings.Unbounded.Unbounded_String);
procedure Get_Line          11/2
  (Item : out Strings.Unbounded.Unbounded_String);
end Ada.Text_IO.Unbounded_IO; 12/2

For an item of type Unbounded_String, the following subprograms are provided: 13/2

procedure Put              14/2
  (File : in File_Type;
   Item : in Strings.Unbounded.Unbounded_String);
  Equivalent to Text_IO.Put (File, Strings.Unbounded.To_String(Item)); 15/2

procedure Put              16/2
  (Item : in Strings.Unbounded.Unbounded_String);
  Equivalent to Text_IO.Put (Strings.Unbounded.To_String(Item)); 17/2

procedure Put_Line          18/2
  (File : in File_Type;
   Item : in Strings.Unbounded.Unbounded_String);
  Equivalent to Text_IO.Put_Line (File, Strings.Unbounded.To_String(Item)); 19/2

procedure Put_Line          20/2
  (Item : in Strings.Unbounded.Unbounded_String);
  Equivalent to Text_IO.Put_Line (Strings.Unbounded.To_String(Item)); 21/2

function Get_Line           22/2
  (File : in File_Type)
  return Strings.Unbounded.Unbounded_String;
  Returns Strings.Unbounded.To_Unbounded_String(Text_IO.Get_Line(File)); 23/2

function Get_Line           24/2
  return Strings.Unbounded.Unbounded_String;
  Returns Strings.Unbounded.To_Unbounded_String(Text_IO.Get_Line); 25/2

procedure Get_Line          26/2
  (File : in File_Type; Item : out Strings.Unbounded.Unbounded_String);
  Equivalent to Item := Get_Line (File); 27/2

procedure Get_Line           28/2
  (Item : out Strings.Unbounded.Unbounded_String);
  Equivalent to Item := Get_Line; 29/2

```

A.11 Wide Text Input-Output and Wide Wide Text Input-Output

The packages `Wide_Text_IO` and `Wide_Wide_Text_IO` provide facilities for input and output in human-readable form. Each file is read or written sequentially, as a sequence of wide characters (or wide wide characters) grouped into lines, and as a sequence of lines grouped into pages.

Static Semantics

The specification of package `Wide_Text_IO` is the same as that for `Text_IO`, except that in each `Get`, `Look_Ahead`, `Get_Immediate`, `Get_Line`, `Put`, and `Put_Line` subprogram, any occurrence of `Character` is replaced by `Wide_Character`, and any occurrence of `String` is replaced by `Wide_String`. Nongeneric equivalents of `Wide_Text_IO.Integer_IO` and `Wide_Text_IO.Float_IO` are provided (as for `Text_IO`) for each predefined numeric type, with names such as `Ada.Integer_Wide_Text_IO`, `Ada.Long_Integer_Wide_Text_IO`, `Ada.Float_Wide_Text_IO`, `Ada.Long_Float_Wide_Text_IO`.

- 3/2 The specification of package `Wide_Wide_Text_IO` is the same as that for `Text_IO`, except that in each `Get`, `Look_Ahead`, `Get_Immediate`, `Get_Line`, `Put`, and `Put_Line` subprogram, any occurrence of `Character` is replaced by `Wide_Wide_Character`, and any occurrence of `String` is replaced by `Wide_Wide_String`. Nongeneric equivalents of `Wide_Wide_Text_IO.Integer_IO` and `Wide_Wide_Text_IO.Float_IO` are provided (as for `Text_IO`) for each predefined numeric type, with names such as `Ada.Integer_Wide_Wide_Text_IO`, `Ada.Long_Integer_Wide_Wide_Text_IO`, `Ada.Float_Wide_Wide_Text_IO`, `Ada.Long_Float_Wide_Wide_Text_IO`.
- 4/2 The specification of package `Wide_Text_IO.Wide_Bounded_IO` is the same as that for `Text_IO.Bounded_IO`, except that any occurrence of `Bounded_String` is replaced by `Wide_Bounded_String`, and any occurrence of package `Bounded` is replaced by `Wide_Bounded`. The specification of package `Wide_Wide_Text_IO.Wide_Wide_Bounded_IO` is the same as that for `Text_IO.Bounded_IO`, except that any occurrence of `Bounded_String` is replaced by `Wide_Wide_Bounded_String`, and any occurrence of package `Bounded` is replaced by `Wide_Wide_Bounded`.
- 5/2 The specification of package `Wide_Text_IO.Wide_Unbounded_IO` is the same as that for `Text_IO.Unbounded_IO`, except that any occurrence of `Unbounded_String` is replaced by `Wide_Unbounded_String`, and any occurrence of package `Unbounded` is replaced by `Wide_Unbounded`. The specification of package `Wide_Wide_Text_IO.Wide_Wide_Unbounded_IO` is the same as that for `Text_IO.Unbounded_IO`, except that any occurrence of `Unbounded_String` is replaced by `Wide_Wide_Unbounded_String`, and any occurrence of package `Unbounded` is replaced by `Wide_Wide_Unbounded`.

A.12 Stream Input-Output

- 1/2 The packages `Streams.Stream_IO`, `Text_IO.Text_Streams`, `Wide_Text_IO.Text_Streams`, and `Wide_Wide_Text_IO.Text_Streams` provide stream-oriented operations on files.

A.12.1 The Package Streams.Stream_IO

- 1 The subprograms in the child package `Streams.Stream_IO` provide control over stream files. Access to a stream file is either sequential, via a call on `Read` or `Write` to transfer an array of stream elements, or positional (if supported by the implementation for the given file), by specifying a relative index for an element. Since a stream file can be converted to a `Stream_Access` value, calling stream-oriented attribute subprograms of different element types with the same `Stream_Access` value provides heterogeneous input-output. See 13.13 for a general discussion of streams.

Static Semantics

- 1.1/1 The elements of a stream file are stream elements. If positioning is supported for the specified external file, a current index and current size are maintained for the file as described in A.8. If positioning is not supported, a current index is not maintained, and the current size is implementation defined.

- 2 The library package `Streams.Stream_IO` has the following declaration:

```

3   with Ada.IO_Exceptions;
4   package Ada.Streams.Stream_IO is
5     type Stream_Access is access all Root_Stream_Type'Class;
6     type File_Type is limited private;
7     type File_Mode is (In_File, Out_File, Append_File);
8     type Count      is range 0 .. implementation-defined;
9     subtype Positive_Count is Count range 1 .. Count'Last;
-- Index into file, in stream elements.

```

```

procedure Create (File : in out File_Type;                                8
                  Mode : in File_Mode := Out_File;
                  Name : in String   := "";
                  Form : in String   := "") ;
procedure Open  (File : in out File_Type;                                9
                  Mode : in File_Mode;
                  Name : in String;
                  Form : in String := "") ;
procedure Close  (File : in out File_Type);                            10
procedure Delete (File : in out File_Type);
procedure Reset  (File : in out File_Type; Mode : in File_Mode);
procedure Reset  (File : in out File_Type);

function Mode (File : in File_Type) return File_Mode;                 11
function Name (File : in File_Type) return String;
function Form (File : in File_Type) return String;

function Is_Open    (File : in File_Type) return Boolean;              12
function End_Of_File (File : in File_Type) return Boolean;
function Stream (File : in File_Type) return Stream_Access;           13
--  Return stream access for use with TInput and TOutput

This paragraph was deleted.                                         14/1

--  Read array of stream elements from file                           15
procedure Read (File : in File_Type;
                Item : out Stream_Element_Array;
                Last : out Stream_Element_Offset;
                From : in Positive_Count);

procedure Read (File : in File_Type;                                16
                Item : out Stream_Element_Array;
                Last : out Stream_Element_Offset);

This paragraph was deleted.                                         17/1

--  Write array of stream elements into file                         18
procedure Write (File : in File_Type;
                 Item : in Stream_Element_Array;
                 To   : in Positive_Count);

procedure Write (File : in File_Type;                                19
                 Item : in Stream_Element_Array);

This paragraph was deleted.                                         20/1

--  Operations on position within file
procedure Set_Index(File : in File_Type; To : in Positive_Count);  21
function Index(File : in File_Type) return Positive_Count;        22
function Size (File : in File_Type) return Count;                  23
procedure Set_Mode(File : in out File_Type; Mode : in File_Mode);  24
procedure Flush(File : in File_Type);                             25/1

--  exceptions
Status_Error : exception renames IO_Exceptions.Status_Error;
Mode_Error   : exception renames IO_Exceptions.Mode_Error;
Name_Error   : exception renames IO_Exceptions.Name_Error;
Use_Error    : exception renames IO_Exceptions.Use_Error;
Device_Error : exception renames IO_Exceptions.Device_Error;
End_Error    : exception renames IO_Exceptions.End_Error;
Data_Error   : exception renames IO_Exceptions.Data_Error;

private
  ... -- not specified by the language
end Ada.Streams.Stream_IO;                                         27

```

The type File_Type needs finalization (see 7.6). 27.1/2

- 28/2 The subprograms given in subclause A.8.2 for the control of external files (Create, Open, Close, Delete, Reset, Mode, Name, Form, and Is_Open) are available for stream files.
- 28.1/2 The End_Of_File function:
- 28.2/2 • Propagates Mode_Error if the mode of the file is not In_File;
 - 28.3/2 • If positioning is supported for the given external file, the function returns True if the current index exceeds the size of the external file; otherwise it returns False;
 - 28.4/2 • If positioning is not supported for the given external file, the function returns True if no more elements can be read from the given file; otherwise it returns False.
- 28.5/2 The Set_Mode procedure sets the mode of the file. If the new mode is Append_File, the file is positioned to its end; otherwise, the position in the file is unchanged.
- 28.6/1 The Flush procedure synchronizes the external file with the internal file (by flushing any internal buffers) without closing the file or changing the position. Mode_Error is propagated if the mode of the file is In_File.
- 29/1 The Stream function returns a Stream_Access result from a File_Type object, thus allowing the stream-oriented attributes Read, Write, Input, and Output to be used on the same file for multiple types. Stream propagates Status_Error if File is not open.
- 30/2 The procedures Read and Write are equivalent to the corresponding operations in the package Streams. Read propagates Mode_Error if the mode of File is not In_File. Write propagates Mode_Error if the mode of File is not Out_File or Append_File. The Read procedure with a Positive_Count parameter starts reading at the specified index. The Write procedure with a Positive_Count parameter starts writing at the specified index. For a file that supports positioning, Read without a Positive_Count parameter starts reading at the current index, and Write without a Positive_Count parameter starts writing at the current index.
- 30.1/1 The Size function returns the current size of the file.
- 31/1 The Index function returns the current index.
- 32 The Set_Index procedure sets the current index to the specified value.
- 32.1/1 If positioning is supported for the external file, the current index is maintained as follows:
- 32.2/1 • For Open and Create, if the Mode parameter is Append_File, the current index is set to the current size of the file plus one; otherwise, the current index is set to one.
- 32.3/1 • For Reset, if the Mode parameter is Append_File, or no Mode parameter is given and the current mode is Append_File, the current index is set to the current size of the file plus one; otherwise, the current index is set to one.
- 32.4/1 • For Set_Mode, if the new mode is Append_File, the current index is set to current size plus one; otherwise, the current index is unchanged.
- 32.5/1 • For Read and Write without a Positive_Count parameter, the current index is incremented by the number of stream elements read or written.
- 32.6/1 • For Read and Write with a Positive_Count parameter, the value of the current index is set to the value of the Positive_Count parameter plus the number of stream elements read or written.
- 33 If positioning is not supported for the given file, then a call of Index or Set_Index propagates Use_Error. Similarly, a call of Read or Write with a Positive_Count parameter propagates Use_Error.

Paragraphs 34 through 36 were deleted.

Erroneous Execution

If the File_Type object passed to the Stream function is later closed or finalized, and the stream-oriented attributes are subsequently called (explicitly or implicitly) on the Stream_Access value returned by Stream, execution is erroneous. This rule applies even if the File_Type object was opened again after it had been closed.

36.1/1

A.12.2 The Package Text_IO.Text_Streams

The package Text_IO.Text_Streams provides a function for treating a text file as a stream.

1

Static Semantics

The library package Text_IO.Text_Streams has the following declaration:

```
with Ada.Streams;
package Ada.Text_IO.Text_Streams is
    type Stream_Access is access all Streams.Root_Stream_Type'Class;
    function Stream (File : in File_Type) return Stream_Access;
end Ada.Text_IO.Text_Streams;
```

2

3

4

The Stream function has the same effect as the corresponding function in Streams.Stream_IO.

5

NOTES

34 The ability to obtain a stream for a text file allows Current_Input, Current_Output, and Current_Error to be processed with the functionality of streams, including the mixing of text and binary input-output, and the mixing of binary input-output for different types.

6

35 Performing operations on the stream associated with a text file does not affect the column, line, or page counts.

7

A.12.3 The Package Wide_Text_IO.Text_Streams

The package Wide_Text_IO.Text_Streams provides a function for treating a wide text file as a stream.

1

Static Semantics

The library package Wide_Text_IO.Text_Streams has the following declaration:

```
with Ada.Streams;
package Ada.Wide_Text_IO.Text_Streams is
    type Stream_Access is access all Streams.Root_Stream_Type'Class;
    function Stream (File : in File_Type) return Stream_Access;
end Ada.Wide_Text_IO.Text_Streams;
```

2

3

4

The Stream function has the same effect as the corresponding function in Streams.Stream_IO.

5

A.12.4 The Package Wide_Wide_Text_IO.Text_Streams

The package Wide_Wide_Text_IO.Text_Streams provides a function for treating a wide wide text file as a stream.

1/2

Static Semantics

The library package Wide_Wide_Text_IO.Text_Streams has the following declaration:

```
with Ada.Streams;
package Ada.Wide_Wide_Text_IO.Text_Streams is
    type Stream_Access is access all Streams.Root_Stream_Type'Class;
    function Stream (File : in File_Type) return Stream_Access;
end Ada.Wide_Wide_Text_IO.Text_Streams;
```

2/2

3/2

4/2

- 5/2 The Stream function has the same effect as the corresponding function in Streams.Stream_IO.

A.13 Exceptions in Input-Output

- 1 The package IO_Exceptions defines the exceptions needed by the predefined input-output packages.

Static Semantics

- 2 The library package IO_Exceptions has the following declaration:

```

3   package Ada.IO_Exceptions is
4     pragma Pure(IO_Exceptions);
5     Status_Error : exception;
6     Mode_Error   : exception;
7     Name_Error   : exception;
8     Use_Error    : exception;
9     Device_Error : exception;
10    End_Error    : exception;
11    Data_Error   : exception;
12    Layout_Error : exception;
13
14   end Ada.IO_Exceptions;

```

- 6 If more than one error condition exists, the corresponding exception that appears earliest in the following list is the one that is propagated.

- 7 The exception Status_Error is propagated by an attempt to operate upon a file that is not open, and by an attempt to open a file that is already open.

- 8 The exception Mode_Error is propagated by an attempt to read from, or test for the end of, a file whose current mode is Out_File or Append_File, and also by an attempt to write to a file whose current mode is In_File. In the case of Text_IO, the exception Mode_Error is also propagated by specifying a file whose current mode is Out_File or Append_File in a call of Set_Input, Skip_Line, End_Of_Line, Skip_Page, or End_Of_Page; and by specifying a file whose current mode is In_File in a call of Set_Output, Set_Line_Length, Set_Page_Length, Line_Length, Page_Length, New_Line, or New_Page.

- 9 The exception Name_Error is propagated by a call of Create or Open if the string given for the parameter Name does not allow the identification of an external file. For example, this exception is propagated if the string is improper, or, alternatively, if either none or more than one external file corresponds to the string.

- 10 The exception Use_Error is propagated if an operation is attempted that is not possible for reasons that depend on characteristics of the external file. For example, this exception is propagated by the procedure Create, among other circumstances, if the given mode is Out_File but the form specifies an input only device, if the parameter Form specifies invalid access rights, or if an external file with the given name already exists and overwriting is not allowed.

- 11 The exception Device_Error is propagated if an input-output operation cannot be completed because of a malfunction of the underlying system.

- 12 The exception End_Error is propagated by an attempt to skip (read past) the end of a file.

- 13 The exception Data_Error can be propagated by the procedure Read (or by the Read attribute) if the element read cannot be interpreted as a value of the required subtype. This exception is also propagated by a procedure Get (defined in the package Text_IO) if the input character sequence fails to satisfy the required syntax, or if the value input does not belong to the range of the required subtype.

The exception Layout_Error is propagated (in text input-output) by Col, Line, or Page if the value returned exceeds Count'Last. The exception Layout_Error is also propagated on output by an attempt to set column or line numbers in excess of specified maximum line or page lengths, respectively (excluding the unbounded cases). It is also propagated by an attempt to Put too many characters to a string.

Documentation Requirements

The implementation shall document the conditions under which Name_Error, Use_Error and Device_Error are propagated.

Implementation Permissions

If the associated check is too complex, an implementation need not propagate Data_Error as part of a procedure Read (or the Read attribute) if the value read cannot be interpreted as a value of the required subtype.

Erroneous Execution

If the element read by the procedure Read (or by the Read attribute) cannot be interpreted as a value of the required subtype, but this is not detected and Data_Error is not propagated, then the resulting value can be abnormal, and subsequent references to the value can lead to erroneous execution, as explained in 13.9.1.

A.14 File Sharing

Dynamic Semantics

It is not specified by the language whether the same external file can be associated with more than one file object. If such sharing is supported by the implementation, the following effects are defined:

- Operations on one text file object do not affect the column, line, and page numbers of any other file object.
- *This paragraph was deleted.*
- For direct and stream files, the current index is a property of each file object; an operation on one file object does not affect the current index of any other file object.
- For direct and stream files, the current size of the file is a property of the external file.

All other effects are identical.

A.15 The Package Command_Line

The package Command_Line allows a program to obtain the values of its arguments and to set the exit status code to be returned on normal termination.

Static Semantics

The library package Ada.Command_Line has the following declaration:

```
package Ada.Command_Line is
  pragma Preelaborate(Command_Line);
  function Argument_Count return Natural;
  function Argument (Number : in Positive) return String;
  function Command_Name return String;
  type Exit_Status is implementation-defined integer type;
```

```

8      Success : constant Exit_Status;
9      Failure : constant Exit_Status;
10     procedure Set_Exit_Status (Code : in Exit_Status);
11     private
12       ... -- not specified by the language
13   end Ada.Command_Line;

14   function Argument_Count return Natural;
15   If the external execution environment supports passing arguments to a program, then
16   Argument_Count returns the number of arguments passed to the program invoking the function.
17   Otherwise it returns 0. The meaning of "number of arguments" is implementation defined.

18   function Argument (Number : in Positive) return String;
19   If the external execution environment supports passing arguments to a program, then Argument
20   returns an implementation-defined value corresponding to the argument at relative position
21   Number. If Number is outside the range 1..Argument_Count, then Constraint_Error is
22   propagated.

23   function Command_Name return String;
24   If the external execution environment supports passing arguments to a program, then Command_Name
25   returns an implementation-defined value corresponding to the name of the
26   command invoking the program; otherwise Command_Name returns the null string.

27 type Exit_Status is implementation-defined integer type;
28   The type Exit_Status represents the range of exit status values supported by the external
29   execution environment. The constants Success and Failure correspond to success and failure,
30   respectively.

31   procedure Set_Exit_Status (Code : in Exit_Status);
32   If the external execution environment supports returning an exit status from a program, then
33   Set_Exit_Status sets Code as the status. Normal termination of a program returns as the exit
34   status the value most recently set by Set_Exit_Status, or, if no such value has been set, then the
35   value Success. If a program terminates abnormally, the status set by Set_Exit_Status is ignored,
36   and an implementation-defined exit status value is set.

37   If the external execution environment does not support returning an exit value from a program,
38   then Set_Exit_Status does nothing.

```

Implementation Permissions

21 An alternative declaration is allowed for package Command_Line if different functionality is appropriate
for the external execution environment.

22 NOTES

36 Argument_Count, Argument, and Command_Name correspond to the C language's argc, argv[n] (for n>0) and
argv[0], respectively.

A.16 The Package Directories

The package Directories provides operations for manipulating files and directories, and their names.

1/2

Static Semantics

The library package Directories has the following declaration:

2/2

```

with Ada.IO_Exceptions;
with Ada.Calendar;
package Ada.Directories is
  -- Directory and file operations:
  function Current_Directory return String;                                     3/2
  procedure Set_Directory (Directory : in String);                                5/2
  procedure Create_Directory (New_Directory : in String;                          6/2
                             Form          : in String := "");                         7/2
  procedure Delete_Directory (Directory : in String);                            8/2
  procedure Create_Path (New_Directory : in String;                           9/2
                        Form          : in String := "");                         10/2
  procedure Delete_Tree (Directory : in String);                               11/2
  procedure Delete_File (Name : in String);                                    12/2
  procedure Rename (Old_Name, New_Name : in String);                           13/2
  procedure Copy_File (Source_Name,
                       Target_Name : in String;                         14/2
                       Form          : in String := "");                     15/2
  -- File and directory name operations:
  function Full_Name (Name : in String) return String;                         16/2
  function Simple_Name (Name : in String) return String;                        17/2
  function Containing_Directory (Name : in String) return String;              18/2
  function Extension (Name : in String) return String;                         19/2
  function Base_Name (Name : in String) return String;                         20/2
  function Compose (Containing_Directory : in String := "";
                    Name           : in String;
                    Extension       : in String := "") return String;
  -- File and directory queries:
  type File_Kind is (Directory, Ordinary_File, Special_File);                22/2
  type File_Size is range 0 .. implementation-defined;                         23/2
  function Exists (Name : in String) return Boolean;                           24/2
  function Kind (Name : in String) return File_Kind;                           25/2
  function Size (Name : in String) return File_Size;                           26/2
  function Modification_Time (Name : in String) return Ada.Calendar.Time;      27/2
  -- Directory searching:
  type Directory_Entry_Type is limited private;                                28/2
  type Filter_Type is array (File_Kind) of Boolean;                            29/2
  type Search_Type is limited private;                                         30/2
  procedure Start_Search (Search    : in out Search_Type;
                         Directory : in String;
                         Pattern   : in String;
                         Filter    : in Filter_Type := (others => True));        31/2
  procedure End_Search (Search : in out Search_Type);                         32/2

```

33/2

```

34/2      function More_Entries (Search : in Search_Type) return Boolean;
35/2      procedure Get_Next_Entry (Search : in out Search_Type;
36/2          Directory_Entry : out Directory_Entry_Type);
36/2      procedure Search (
37/2          Directory : in String;
38/2          Pattern   : in String;
39/2          Filter    : in Filter_Type := (others => True);
40/2          Process   : not null access procedure (
41/2              Directory_Entry : in Directory_Entry_Type));
42/2      -- Operations on Directory Entries:
43/2      function Simple_Name (Directory_Entry : in Directory_Entry_Type)
44/2          return String;
45/2      function Full_Name (Directory_Entry : in Directory_Entry_Type)
46/2          return String;
47/2      function Kind (Directory_Entry : in Directory_Entry_Type)
48/2          return File_Kind;
49/2      function Size (Directory_Entry : in Directory_Entry_Type)
50/2          return File_Size;
51/2      function Modification_Time (Directory_Entry : in Directory_Entry_Type)
52/2          return Ada.Calendar.Time;
53/2      Status_Error : exception renames Ada.IO_Exceptions.Status_Error;
54/2      Name_Error   : exception renames Ada.IO_Exceptions.Name_Error;
55/2      Use_Error    : exception renames Ada.IO_Exceptions.Use_Error;
56/2      Device_Error : exception renames Ada.IO_Exceptions.Device_Error;
57/2      private
58/2          -- Not specified by the language.
59/2      end Ada.Directories;

```

45/2 External files may be classified as directories, special files, or ordinary files. A *directory* is an external file that is a container for files on the target system. A *special file* is an external file that cannot be created or read by a predefined Ada input-output package. External files that are not special files or directories are called *ordinary files*.

46/2 A *file name* is a string identifying an external file. Similarly, a *directory name* is a string identifying a directory. The interpretation of file names and directory names is implementation-defined.

47/2 The *full name* of an external file is a full specification of the name of the file. If the external environment allows alternative specifications of the name (for example, abbreviations), the full name should not use such alternatives. A full name typically will include the names of all of the directories that contain the item. The *simple name* of an external file is the name of the item, not including any containing directory names. Unless otherwise specified, a file name or directory name parameter in a call to a predefined Ada input-output subprogram can be a full name, a simple name, or any other form of name supported by the implementation.

48/2 The *default directory* is the directory that is used if a directory or file name is not a full name (that is, when the name does not fully identify all of the containing directories).

49/2 A *directory entry* is a single item in a directory, identifying a single external file (including directories and special files).

50/2 For each function that returns a string, the lower bound of the returned value is 1.

51/2 The following file and directory operations are provided:

function Current_Directory return String;	52/2
Returns the full directory name for the current default directory. The name returned shall be suitable for a future call to Set_Directory. The exception Use_Error is propagated if a default directory is not supported by the external environment.	53/2
procedure Set_Directory (Directory : in String);	54/2
Sets the current default directory. The exception Name_Error is propagated if the string given as Directory does not identify an existing directory. The exception Use_Error is propagated if the external environment does not support making Directory (in the absence of Name_Error) a default directory.	55/2
procedure Create_Directory (New_Directory : in String; Form : in String := "");	56/2
Creates a directory with name New_Directory. The Form parameter can be used to give system-dependent characteristics of the directory; the interpretation of the Form parameter is implementation-defined. A null string for Form specifies the use of the default options of the implementation of the new directory. The exception Name_Error is propagated if the string given as New_Directory does not allow the identification of a directory. The exception Use_Error is propagated if the external environment does not support the creation of a directory with the given name (in the absence of Name_Error) and form.	57/2
procedure Delete_Directory (Directory : in String);	58/2
Deletes an existing empty directory with name Directory. The exception Name_Error is propagated if the string given as Directory does not identify an existing directory. The exception Use_Error is propagated if the external environment does not support the deletion of the directory (or some portion of its contents) with the given name (in the absence of Name_Error).	59/2
procedure Create_Path (New_Directory : in String; Form : in String := "");	60/2
Creates zero or more directories with name New_Directory. Each non-existent directory named by New_Directory is created. For example, on a typical Unix system, Create_Path ("~/usr/me/my"); would create directory "me" in directory "usr", then create directory "my" in directory "me". The Form parameter can be used to give system-dependent characteristics of the directory; the interpretation of the Form parameter is implementation-defined. A null string for Form specifies the use of the default options of the implementation of the new directory. The exception Name_Error is propagated if the string given as New_Directory does not allow the identification of any directory. The exception Use_Error is propagated if the external environment does not support the creation of any directories with the given name (in the absence of Name_Error) and form.	61/2
procedure Delete_Tree (Directory : in String);	62/2
Deletes an existing directory with name Directory. The directory and all of its contents (possibly including other directories) are deleted. The exception Name_Error is propagated if the string given as Directory does not identify an existing directory. The exception Use_Error is propagated if the external environment does not support the deletion of the directory or some portion of its contents with the given name (in the absence of Name_Error). If Use_Error is propagated, it is unspecified whether a portion of the contents of the directory is deleted.	63/2

```

64/2   procedure Delete_File (Name : in String);
65/2     Deletes an existing ordinary or special file with name Name. The exception Name_Error is
     propagated if the string given as Name does not identify an existing ordinary or special external
     file. The exception Use_Error is propagated if the external environment does not support the
     deletion of the file with the given name (in the absence of Name_Error).

66/2   procedure Rename (Old_Name, New_Name : in String);
67/2     Renames an existing external file (including directories) with name Old_Name to New_Name.
     The exception Name_Error is propagated if the string given as Old_Name does not identify an
     existing external file. The exception Use_Error is propagated if the external environment does
     not support the renaming of the file with the given name (in the absence of Name_Error). In
     particular, Use_Error is propagated if a file or directory already exists with name New_Name.

68/2   procedure Copy_File (Source_Name,
                           Target_Name : in String;
                           Form         : in String);
69/2     Copies the contents of the existing external file with name Source_Name to an external file with
     name Target_Name. The resulting external file is a duplicate of the source external file. The
     Form parameter can be used to give system-dependent characteristics of the resulting external
     file; the interpretation of the Form parameter is implementation-defined. Exception Name_Error
     is propagated if the string given as Source_Name does not identify an existing external ordinary
     or special file, or if the string given as Target_Name does not allow the identification of an
     external file. The exception Use_Error is propagated if the external environment does not
     support creating the file with the name given by Target_Name and form given by Form, or
     copying of the file with the name given by Source_Name (in the absence of Name_Error).

70/2   The following file and directory name operations are provided:
71/2     function Full_Name (Name : in String) return String;
72/2       Returns the full name corresponding to the file name specified by Name. The exception
     Name_Error is propagated if the string given as Name does not allow the identification of an
     external file (including directories and special files).

73/2     function Simple_Name (Name : in String) return String;
74/2       Returns the simple name portion of the file name specified by Name. The exception Name_Error
     is propagated if the string given as Name does not allow the identification of an external file
     (including directories and special files).

75/2     function Containing_Directory (Name : in String) return String;
76/2       Returns the name of the containing directory of the external file (including directories) identified
     by Name. (If more than one directory can contain Name, the directory name returned is
     implementation-defined.) The exception Name_Error is propagated if the string given as Name
     does not allow the identification of an external file. The exception Use_Error is propagated if the
     external file does not have a containing directory.

77/2     function Extension (Name : in String) return String;
78/2       Returns the extension name corresponding to Name. The extension name is a portion of a simple
     name (not including any separator characters), typically used to identify the file class. If the
     external environment does not have extension names, then the null string is returned. The

```

exception `Name_Error` is propagated if the string given as `Name` does not allow the identification of an external file.

```
function Base_Name (Name : in String) return String;
```

79/2

Returns the base name corresponding to `Name`. The base name is the remainder of a simple name after removing any extension and extension separators. The exception `Name_Error` is propagated if the string given as `Name` does not allow the identification of an external file (including directories and special files).

80/2

```
function Compose (Containing_Directory : in String := "";  
                  Name           : in String;  
                  Extension       : in String := "") return String;
```

81/2

Returns the name of the external file with the specified `Containing_Directory`, `Name`, and `Extension`. If `Extension` is the null string, then `Name` is interpreted as a simple name; otherwise `Name` is interpreted as a base name. The exception `Name_Error` is propagated if the string given as `Containing_Directory` is not null and does not allow the identification of a directory, or if the string given as `Extension` is not null and is not a possible extension, or if the string given as `Name` is not a possible simple name (if `Extension` is null) or base name (if `Extension` is non-null).

82/2

The following file and directory queries and types are provided:

83/2

```
type File_Kind is (Directory, Ordinary_File, Special_File);
```

84/2

The type `File_Kind` represents the kind of file represented by an external file or directory.

85/2

```
type File_Size is range 0 .. implementation-defined;
```

86/2

The type `File_Size` represents the size of an external file.

87/2

```
function Exists (Name : in String) return Boolean;
```

88/2

Returns True if an external file represented by `Name` exists, and False otherwise. The exception `Name_Error` is propagated if the string given as `Name` does not allow the identification of an external file (including directories and special files).

89/2

```
function Kind (Name : in String) return File_Kind;
```

90/2

Returns the kind of external file represented by `Name`. The exception `Name_Error` is propagated if the string given as `Name` does not allow the identification of an existing external file.

91/2

```
function Size (Name : in String) return File_Size;
```

92/2

Returns the size of the external file represented by `Name`. The size of an external file is the number of stream elements contained in the file. If the external file is not an ordinary file, the result is implementation-defined. The exception `Name_Error` is propagated if the string given as `Name` does not allow the identification of an existing external file. The exception `Constraint_Error` is propagated if the file size is not a value of type `File_Size`.

93/2

```
function Modification_Time (Name : in String) return Ada.Calendar.Time;
```

94/2

Returns the time that the external file represented by `Name` was most recently modified. If the external file is not an ordinary file, the result is implementation-defined. The exception `Name_Error` is propagated if the string given as `Name` does not allow the identification of an existing external file. The exception `Use_Error` is propagated if the external environment does not support reading the modification time of the file with the name given by `Name` (in the absence of `Name_Error`).

95/2

96/2 The following directory searching operations and types are provided:

97/2 **type** Directory_Entry_Type **is limited private;**

98/2 The type Directory_Entry_Type represents a single item in a directory. These items can only be created by the Get_Next_Entry procedure in this package. Information about the item can be obtained from the functions declared in this package. A default-initialized object of this type is invalid; objects returned from Get_Next_Entry are valid.

99/2 **type** Filter_Type **is array (File_Kind) of Boolean;**

100/2 The type Filter_Type specifies which directory entries are provided from a search operation. If the Directory component is True, directory entries representing directories are provided. If the Ordinary_File component is True, directory entries representing ordinary files are provided. If the Special_File component is True, directory entries representing special files are provided.

101/2 **type** Search_Type **is limited private;**

102/2 The type Search_Type contains the state of a directory search. A default-initialized Search_Type object has no entries available (function MoreEntries returns False). Type Search_Type needs finalization (see 7.6).

103/2 **procedure** Start_Search (Search : **in out** Search_Type;
 Directory : **in** String;
 Pattern : **in** String;
 Filter : **in** Filter_Type := (**others** => True));

104/2 Starts a search in the directory named by Directory for entries matching Pattern. Pattern represents a pattern for matching file names. If Pattern is null, all items in the directory are matched; otherwise, the interpretation of Pattern is implementation-defined. Only items that match Filter will be returned. After a successful call on Start_Search, the object Search may have entries available, but it may have no entries available if no files or directories match Pattern and Filter. The exception Name_Error is propagated if the string given by Directory does not identify an existing directory, or if Pattern does not allow the identification of any possible external file or directory. The exception Use_Error is propagated if the external environment does not support the searching of the directory with the given name (in the absence of Name_Error). When Start_Search propagates Name_Error or Use_Error, the object Search will have no entries available.

105/2 **procedure** End_Search (Search : **in out** Search_Type);

106/2 Ends the search represented by Search. After a successful call on End_Search, the object Search will have no entries available.

107/2 **function** MoreEntries (Search : **in** Search_Type) **return** Boolean;

108/2 Returns True if more entries are available to be returned by a call to Get_Next_Entry for the specified search object, and False otherwise.

109/2 **procedure** Get_Next_Entry (Search : **in out** Search_Type;
 Directory_Entry : **out** Directory_Entry_Type);

110/2 Returns the next Directory_Entry for the search described by Search that matches the pattern and filter. If no further matches are available, Status_Error is raised. It is implementation-defined as to whether the results returned by this routine are altered if the contents of the directory are altered while the Search object is valid (for example, by another program). The exception Use_Error is propagated if the external environment does not support continued searching of the directory represented by Search.

```
procedure Search (
  Directory : in String;
  Pattern   : in String;
  Filter    : in Filter_Type := (others => True);
  Process   : not null access procedure (
    Directory_Entry : in Directory_Entry_Type));
```

111/2

Searches in the directory named by Directory for entries matching Pattern. The subprogram designated by Process is called with each matching entry in turn. Pattern represents a pattern for matching file names. If Pattern is null, all items in the directory are matched; otherwise, the interpretation of Pattern is implementation-defined. Only items that match Filter will be returned. The exception Name_Error is propagated if the string given by Directory does not identify an existing directory, or if Pattern does not allow the identification of any possible external file or directory. The exception Use_Error is propagated if the external environment does not support the searching of the directory with the given name (in the absence of Name_Error).

112/2

```
function Simple_Name (Directory_Entry : in Directory_Entry_Type)
  return String;
```

113/2

Returns the simple external name of the external file (including directories) represented by Directory_Entry. The format of the name returned is implementation-defined. The exception Status_Error is propagated if Directory_Entry is invalid.

114/2

```
function Full_Name (Directory_Entry : in Directory_Entry_Type)
  return String;
```

115/2

Returns the full external name of the external file (including directories) represented by Directory_Entry. The format of the name returned is implementation-defined. The exception Status_Error is propagated if Directory_Entry is invalid.

116/2

```
function Kind (Directory_Entry : in Directory_Entry_Type)
  return File_Kind;
```

117/2

Returns the kind of external file represented by Directory_Entry. The exception Status_Error is propagated if Directory_Entry is invalid.

118/2

```
function Size (Directory_Entry : in Directory_Entry_Type)
  return File_Size;
```

119/2

Returns the size of the external file represented by Directory_Entry. The size of an external file is the number of stream elements contained in the file. If the external file represented by Directory_Entry is not an ordinary file, the result is implementation-defined. The exception Status_Error is propagated if Directory_Entry is invalid. The exception Constraint_Error is propagated if the file size is not a value of type File_Size.

120/2

```
function Modification_Time (Directory_Entry : in Directory_Entry_Type)
  return Ada.Calendar.Time;
```

121/2

Returns the time that the external file represented by Directory_Entry was most recently modified. If the external file represented by Directory_Entry is not an ordinary file, the result is implementation-defined. The exception Status_Error is propagated if Directory_Entry is invalid. The exception Use_Error is propagated if the external environment does not support reading the modification time of the file represented by Directory_Entry.

122/2

Implementation Requirements

- 123/2 For `Copy_File`, if `Source_Name` identifies an existing external ordinary file created by a predefined Ada input-output package, and `Target_Name` and `Form` can be used in the `Create` operation of that input-output package with mode `Out_File` without raising an exception, then `Copy_File` shall not propagate `Use_Error`.

Implementation Advice

- 124/2 If other information about a file (such as the owner or creation date) is available in a directory entry, the implementation should provide functions in a child package `Directories`.Information to retrieve it.
- 125/2 `Start_Search` and `Search` should raise `Use_Error` if `Pattern` is malformed, but not if it could represent a file in the directory but does not actually do so.
- 126/2 `Rename` should be supported at least when both `New_Name` and `Old_Name` are simple names and `New_Name` does not identify an existing external file.

NOTES

- 127/2 37 The operations `Containing_Directory`, `Full_Name`, `Simple_Name`, `Base_Name`, `Extension`, and `Compose` operate on file names, not external files. The files identified by these operations do not need to exist. `Name_Error` is raised only if the file name is malformed and cannot possibly identify a file. Of these operations, only the result of `Full_Name` depends on the current default directory; the result of the others depends only on their parameters.
- 128/2 38 Using access types, values of `Search_Type` and `Directory_Entry_Type` can be saved and queried later. However, another task or application can modify or delete the file represented by a `Directory_Entry_Type` value or the directory represented by a `Search_Type` value; such a value can only give the information valid at the time it is created. Therefore, long-term storage of these values is not recommended.
- 129/2 39 If the target system does not support directories inside of directories, then `Kind` will never return `Directory` and `Containing_Directory` will always raise `Use_Error`.
- 130/2 40 If the target system does not support creation or deletion of directories, then `Create_Directory`, `Create_Path`, `Delete_Directory`, and `Delete_Tree` will always propagate `Use_Error`.
- 131/2 41 To move a file or directory to a different location, use `Rename`. Most target systems will allow renaming of files from one directory to another. If the target file or directory might already exist, it should be deleted first.

A.17 The Package Environment_Variables

- 1/2 The package `Environment_Variables` allows a program to read or modify environment variables. Environment variables are name-value pairs, where both the name and value are strings. The definition of what constitutes an *environment variable*, and the meaning of the name and value, are implementation defined.

Static Semantics

- 2/2 The library package `Environment_Variables` has the following declaration:

```

3/2   package Ada.Environment_Variables is
4/2     pragma Preelaborate(Enviroment_Variables);
5/2     function Value (Name : in String) return String;
6/2     function Exists (Name : in String) return Boolean;
7/2     procedure Set (Name : in String; Value : in String);
8/2     procedure Clear (Name : in String);
9/2     procedure Clear;
      procedure Iterate (
        Process : not null access procedure (Name, Value : in String));
end Ada.Environment_Variables;
```

function Value (Name : in String) return String;	10/2
If the external execution environment supports environment variables, then Value returns the value of the environment variable with the given name. If no environment variable with the given name exists, then Constraint_Error is propagated. If the execution environment does not support environment variables, then Program_Error is propagated.	11/2
function Exists (Name : in String) return Boolean;	12/2
If the external execution environment supports environment variables and an environment variable with the given name currently exists, then Exists returns True; otherwise it returns False.	13/2
procedure Set (Name : in String; Value : in String);	14/2
If the external execution environment supports environment variables, then Set first clears any existing environment variable with the given name, and then defines a single new environment variable with the given name and value. Otherwise Program_Error is propagated.	15/2
If implementation-defined circumstances prohibit the definition of an environment variable with the given name and value, then Constraint_Error is propagated.	16/2
It is implementation defined whether there exist values for which the call Set(Name, Value) has the same effect as Clear (Name).	17/2
procedure Clear (Name : in String);	18/2
If the external execution environment supports environment variables, then Clear deletes all existing environment variable with the given name. Otherwise Program_Error is propagated.	19/2
procedure Clear;	20/2
If the external execution environment supports environment variables, then Clear deletes all existing environment variables. Otherwise Program_Error is propagated.	21/2
procedure Iterate (Process : not null access procedure (Name, Value : in String));	22/2
If the external execution environment supports environment variables, then Iterate calls the subprogram designated by Process for each existing environment variable, passing the name and value of that environment variable. Otherwise Program_Error is propagated.	23/2
If several environment variables exist that have the same name, Process is called once for each such variable.	24/2

Bounded (Run-Time) Errors

It is a bounded error to call Value if more than one environment variable exists with the given name; the possible outcomes are that:

- one of the values is returned, and that same value is returned in subsequent calls in the absence of changes to the environment; or
- Program_Error is propagated.

Erroneous Execution

Making calls to the procedures Set or Clear concurrently with calls to any subprogram of package Environment_Variables, or to any instantiation of Iterate, results in erroneous execution.

- 29/2 Making calls to the procedures Set or Clear in the actual subprogram corresponding to the Process parameter of Iterate results in erroneous execution.

Documentation Requirements

- 30/2 An implementation shall document how the operations of this package behave if environment variables are changed by external mechanisms (for instance, calling operating system services).

Implementation Permissions

- 31/2 An implementation running on a system that does not support environment variables is permitted to define the operations of package Environment_Variables with the semantics corresponding to the case where the external execution environment does support environment variables. In this case, it shall provide a mechanism to initialize a nonempty set of environment variables prior to the execution of a partition.

Implementation Advice

- 32/2 If the execution environment supports subprocesses, the currently defined environment variables should be used to initialize the environment variables of a subprocess.
- 33/2 Changes to the environment variables made outside the control of this package should be reflected immediately in the effect of the operations of this package. Changes to the environment variables made using this package should be reflected immediately in the external execution environment. This package should not perform any buffering of the environment variables.

A.18 Containers

This clause presents the specifications of the package Containers and several child packages, which provide facilities for storing collections of elements.

A variety of sequence and associative containers are provided. Each container includes a *cursor* type. A cursor is a reference to an element within a container. Many operations on cursors are common to all of the containers. A cursor referencing an element in a container is considered to be overlapping with the container object itself.

Within this clause we provide Implementation Advice for the desired average or worst case time complexity of certain operations on a container. This advice is expressed using the Landau symbol $O(X)$. Presuming f is some function of a length parameter N and $t(N)$ is the time the operation takes (on average or worst case, as specified) for the length N , a complexity of $O(f(N))$ means that there exists a finite A such that for any N , $t(N)/f(N) < A$.

If the advice suggests that the complexity should be less than $O(f(N))$, then for any arbitrarily small positive real D , there should exist a positive integer M such that for all $N > M$, $t(N)/f(N) < D$.

A.18.1 The Package Containers

The package Containers is the root of the containers subsystem.

Static Semantics

The library package Containers has the following declaration:

```
package Ada.Containers is
  pragma Pure(Containers);
  type Hash_Type is mod implementation-defined;
  type Count_Type is range 0 .. implementation-defined;
end Ada.Containers;
```

Hash_Type represents the range of the result of a hash function. Count_Type represents the (potential or actual) number of elements of a container.

Implementation Advice

Hash_Type'Modulus should be at least 2^{**32} . Count_Type'Last should be at least $2^{**31}-1$.

A.18.2 The Package Containers.Vectors

The language-defined generic package Containers.Vectors provides private types Vector and Cursor, and a set of operations for each type. A vector container allows insertion and deletion at any position, but it is specifically optimized for insertion and deletion at the high end (the end with the higher index) of the container. A vector container also provides random access to its elements.

A vector container behaves conceptually as an array that expands as necessary as items are inserted. The *length* of a vector is the number of elements that the vector contains. The *capacity* of a vector is the maximum number of elements that can be inserted into the vector prior to it being automatically expanded.

Elements in a vector container can be referred to by an index value of a generic formal type. The first element of a vector always has its index value equal to the lower bound of the formal type.

- 4/2 A vector container may contain *empty elements*. Empty elements do not have a specified value.

Static Semantics

- 5/2 The generic library package Containers.Vectors has the following declaration:

```

6/2   generic
      type Index_Type is range <>;
      type Element_Type is private;
      with function "=" (Left, Right : Element_Type)
           return Boolean is <>;
  package Ada.Containers.Vectors is
    pragma Preelaborate(Vectors);
7/2   subtype Extended_Index is
      Index_Type'Base range
        Index_Type'First-1 ..
        Index_Type'Min (Index_Type'Base'Last - 1, Index_Type'Last) + 1;
      No_Index : constant Extended_Index := Extended_Index'First;
8/2   type Vector is tagged private;
  pragma Preelaborable_Initialization(Vector);
9/2   type Cursor is private;
  pragma Preelaborable_Initialization(Cursor);
10/2  Empty_Vector : constant Vector;
11/2  No_Element : constant Cursor;
12/2  function "=" (Left, Right : Vector) return Boolean;
13/2  function To_Vector (Length : Count_Type) return Vector;
14/2  function To_Vector
      (New_Item : Element_Type;
       Length   : Count_Type) return Vector;
15/2  function "&" (Left, Right : Vector) return Vector;
16/2  function "&" (Left   : Vector;
                  Right  : Element_Type) return Vector;
17/2  function "&" (Left   : Element_Type;
                  Right  : Vector) return Vector;
18/2  function "&" (Left, Right  : Element_Type) return Vector;
19/2  function Capacity (Container : Vector) return Count_Type;
20/2  procedure Reserve_Capacity (Container : in out Vector;
                                 Capacity   : in      Count_Type);
21/2  function Length (Container : Vector) return Count_Type;
22/2  procedure Set_Length (Container : in out Vector;
                           Length    : in      Count_Type);
23/2  function Is_Empty (Container : Vector) return Boolean;
24/2  procedure Clear (Container : in out Vector);
25/2  function To_Cursor (Container : Vector;
                         Index     : Extended_Index) return Cursor;
26/2  function To_Index (Position  : Cursor) return Extended_Index;
27/2  function Element (Container : Vector;
                         Index     : Index_Type)
           return Element_Type;
28/2  function Element (Position : Cursor) return Element_Type;
29/2  procedure Replace_Element (Container : in out Vector;
                               Index     : in      Index_Type;
                               New_Item : in      Element_Type);

```

procedure Replace_Element (Container : in out Vector;	30/2
Position : in Cursor;	
New_item : in Element_Type);	
procedure Query_Element	31/2
(Container : in Vector;	
Index : in Index_Type;	
Process : not null access procedure (Element : in Element_Type));	
procedure Query_Element	32/2
(Position : in Cursor;	
Process : not null access procedure (Element : in Element_Type));	
procedure Update_Element	33/2
(Container : in out Vector;	
Index : in Index_Type;	
Process : not null access procedure	
(Element : in out Element_Type));	
procedure Update_Element	34/2
(Container : in out Vector;	
Position : in Cursor;	
Process : not null access procedure	
(Element : in out Element_Type));	
procedure Move (Target : in out Vector;	35/2
Source : in out Vector);	
procedure Insert (Container : in out Vector;	36/2
Before : in Extended_Index;	
New_Item : in Vector);	
procedure Insert (Container : in out Vector;	37/2
Before : in Cursor;	
New_Item : in Vector);	
procedure Insert (Container : in out Vector;	38/2
Before : in Cursor;	
New_Item : in Vector;	
Position : out Cursor);	
procedure Insert (Container : in out Vector;	39/2
Before : in Extended_Index;	
New_Item : in Element_Type;	
Count : in Count_Type := 1);	
procedure Insert (Container : in out Vector;	40/2
Before : in Cursor;	
New_Item : in Element_Type;	
Count : in Count_Type := 1);	
procedure Insert (Container : in out Vector;	41/2
Before : in Cursor;	
New_Item : in Element_Type;	
Position : out Cursor;	
Count : in Count_Type := 1);	
procedure Insert (Container : in out Vector;	42/2
Before : in Extended_Index;	
Count : in Count_Type := 1);	
procedure Insert (Container : in out Vector;	43/2
Before : in Cursor;	
Position : out Cursor;	
Count : in Count_Type := 1);	
procedure Prepend (Container : in out Vector;	44/2
New_Item : in Vector);	
procedure Prepend (Container : in out Vector;	45/2
New_Item : in Element_Type;	
Count : in Count_Type := 1);	
procedure Append (Container : in out Vector;	46/2
New_Item : in Vector);	

```

47/2      procedure Append (Container : in out Vector;
                           New_Item  : in      Element_Type;
                           Count      : in      Count_Type := 1);
48/2      procedure Insert_Space (Container : in out Vector;
                                 Before    : in      Extended_Index;
                                 Count     : in      Count_Type := 1);
49/2      procedure Insert_Space (Container : in out Vector;
                                 Before    : in      Cursor;
                                 Position   : out    Cursor;
                                 Count     : in      Count_Type := 1);
50/2      procedure Delete (Container : in out Vector;
                           Index    : in      Extended_Index;
                           Count     : in      Count_Type := 1);
51/2      procedure Delete (Container : in out Vector;
                           Position  : in out Cursor;
                           Count     : in      Count_Type := 1);
52/2      procedure Delete_First (Container : in out Vector;
                                 Count     : in      Count_Type := 1);
53/2      procedure Delete_Last (Container : in out Vector;
                               Count     : in      Count_Type := 1);
54/2      procedure Reverse_Elements (Container : in out Vector);
55/2      procedure Swap (Container : in out Vector;
                           I, J      : in      Index_Type);
56/2      procedure Swap (Container : in out Vector;
                           I, J      : in      Cursor);
57/2      function First_Index (Container : Vector) return Index_Type;
58/2      function First (Container : Vector) return Cursor;
59/2      function First_Element (Container : Vector)
                           return Element_Type;
60/2      function Last_Index (Container : Vector) return Extended_Index;
61/2      function Last (Container : Vector) return Cursor;
62/2      function Last_Element (Container : Vector)
                           return Element_Type;
63/2      function Next (Position : Cursor) return Cursor;
64/2      procedure Next (Position : in out Cursor);
65/2      function Previous (Position : Cursor) return Cursor;
66/2      procedure Previous (Position : in out Cursor);
67/2      function Find_Index (Container : Vector;
                           Item       : Element_Type;
                           Index      : Index_Type := Index_Type'First)
                           return Extended_Index;
68/2      function Find (Container : Vector;
                           Item       : Element_Type;
                           Position   : Cursor := No_Element)
                           return Cursor;
69/2      function Reverse_Find_Index (Container : Vector;
                                       Item       : Element_Type;
                                       Index      : Index_Type := Index_Type'Last)
                           return Extended_Index;
70/2      function Reverse_Find (Container : Vector;
                               Item       : Element_Type;
                               Position   : Cursor := No_Element)
                           return Cursor;
71/2      function Contains (Container : Vector;
                           Item       : Element_Type) return Boolean;

```

```

function Has_Element (Position : Cursor) return Boolean;           72/2
procedure Iterate
  (Container : in Vector;
   Process   : not null access procedure (Position : in Cursor));
procedure Reverse_Iterate
  (Container : in Vector;
   Process   : not null access procedure (Position : in Cursor));
generic
  with function "<" (Left, Right : Element_Type)
    return Boolean is <>;
package Generic_Sorting is
  function Is_Sorted (Container : Vector) return Boolean;          76/2
  procedure Sort (Container : in out Vector);                      77/2
  procedure Merge (Target   : in out Vector;
                  Source   : in out Vector);                         78/2
end Generic_Sorting;                                              79/2
private
  ... -- not specified by the language                           80/2
end Ada.Containers.Vectors;                                         81/2
82/2
83/2

```

The actual function for the generic formal function "`=`" on `Element_Type` values is expected to define a reflexive and symmetric relationship and return the same result value each time it is called with a particular pair of values. If it behaves in some other manner, the functions defined to use it return an unspecified value. The exact arguments and number of calls of this generic formal function by the functions defined to use it are unspecified.

The type `Vector` is used to represent vectors. The type `Vector` needs finalization (see 7.6). 84/2

`Empty_Vector` represents the empty vector object. It has a length of 0. If an object of type `Vector` is not otherwise initialized, it is initialized to the same value as `Empty_Vector`. 85/2

`No_Element` represents a cursor that designates no element. If an object of type `Cursor` is not otherwise initialized, it is initialized to the same value as `No_Element`. 86/2

The predefined "`=`" operator for type `Cursor` returns True if both cursors are `No_Element`, or designate the same element in the same container. 87/2

Execution of the default implementation of the `Input`, `Output`, `Read`, or `Write` attribute of type `Cursor` raises `Program_Error`. 88/2

`No_Index` represents a position that does not correspond to any element. The subtype `Extended_Index` includes the indices covered by `Index_Type` plus the value `No_Index` and, if it exists, the successor to the `Index_Type'Last`. 89/2

Some operations of this generic package have access-to-subprogram parameters. To ensure such operations are well-defined, they guard against certain actions by the designated subprogram. In particular, some operations check for "tampering with cursors" of a container because they depend on the set of elements of the container remaining constant, and others check for "tampering with elements" of a container because they depend on elements of the container not being replaced. 90/2

A subprogram is said to *tamper with cursors* of a vector object *V* if: 91/2

- it inserts or deletes elements of *V*, that is, it calls the `Insert`, `Insert_Space`, `Clear`, `Delete`, or `Set_Length` procedures with *V* as a parameter; or 92/2
- it finalizes *V*; or 93/2

- 94/2 • it calls the Move procedure with V as a parameter.
- 95/2 A subprogram is said to *tamper with elements* of a vector object V if:
- 96/2 • it tampers with cursors of V ; or
- 97/2 • it replaces one or more elements of V , that is, it calls the Replace_Element, Reverse_Elements, or Swap procedures or the Sort or Merge procedures of an instance of Generic_Sorting with V as a parameter.
- 98/2 **function** " $=$ " (Left, Right : Vector) **return** Boolean;
- 99/2 If Left and Right denote the same vector object, then the function returns True. If Left and Right have different lengths, then the function returns False. Otherwise, it compares each element in Left to the corresponding element in Right using the generic formal equality operator. If any such comparison returns False, the function returns False; otherwise it returns True. Any exception raised during evaluation of element equality is propagated.
- 100/2 **function** To_Vector (Length : Count_Type) **return** Vector;
- 101/2 Returns a vector with a length of Length, filled with empty elements.
- 102/2 **function** To_Vector
 (New_Item : Element_Type;
 Length : Count_Type) **return** Vector;
- 103/2 Returns a vector with a length of Length, filled with elements initialized to the value New_Item.
- 104/2 **function** "&" (Left, Right : Vector) **return** Vector;
- 105/2 Returns a vector comprising the elements of Left followed by the elements of Right.
- 106/2 **function** "&" (Left : Vector;
 Right : Element_Type) **return** Vector;
- 107/2 Returns a vector comprising the elements of Left followed by the element Right.
- 108/2 **function** "&" (Left : Element_Type;
 Right : Vector) **return** Vector;
- 109/2 Returns a vector comprising the element Left followed by the elements of Right.
- 110/2 **function** "&" (Left, Right : Element_Type) **return** Vector;
- 111/2 Returns a vector comprising the element Left followed by the element Right.
- 112/2 **function** Capacity (Container : Vector) **return** Count_Type;
- 113/2 Returns the capacity of Container.
- 114/2 **procedure** Reserve_Capacity (Container : **in out** Vector;
 Capacity : **in** Count_Type);
- 115/2 Reserve_Capacity allocates new internal data structures such that the length of the resulting vector can become at least the value Capacity without requiring an additional call to Reserve_Capacity, and is large enough to hold the current length of Container. Reserve_Capacity then copies the elements into the new data structures and deallocates the old data structures. Any exception raised during allocation is propagated and Container is not modified.
- 116/2 **function** Length (Container : Vector) **return** Count_Type;
- 117/2 Returns the number of elements in Container.

procedure Set_Length (Container : in out Vector; Length : in Count_Type);	118/2
If Length is larger than the capacity of Container, Set_Length calls Reserve_Capacity (Container, Length), then sets the length of the Container to Length. If Length is greater than the original length of Container, empty elements are added to Container; otherwise elements are removed from Container.	119/2
function Is_Empty (Container : Vector) return Boolean; Equivalent to Length (Container) = 0.	120/2 121/2
procedure Clear (Container : in out Vector); Removes all the elements from Container. The capacity of Container does not change.	122/2 123/2
function To_Cursor (Container : Vector; Index : Extended_Index) return Cursor;	124/2
If Index is not in the range First_Index (Container) .. Last_Index (Container), then No_Element is returned. Otherwise, a cursor designating the element at position Index in Container is returned.	125/2
function To_Index (Position : Cursor) return Extended_Index;	126/2
If Position is No_Element, No_Index is returned. Otherwise, the index (within its containing vector) of the element designated by Position is returned.	127/2
function Element (Container : Vector; Index : Index_Type) return Element_Type;	128/2
If Index is not in the range First_Index (Container) .. Last_Index (Container), then Constraint_Error is propagated. Otherwise, Element returns the element at position Index.	129/2
function Element (Position : Cursor) return Element_Type;	130/2
If Position equals No_Element, then Constraint_Error is propagated. Otherwise, Element returns the element designated by Position.	131/2
procedure Replace_Element (Container : in out Vector; Index : in Index_Type; New_Item : in Element_Type);	132/2
If Index is not in the range First_Index (Container) .. Last_Index (Container), then Constraint_Error is propagated. Otherwise Replace_Element assigns the value New_Item to the element at position Index. Any exception raised during the assignment is propagated. The element at position Index is not an empty element after successful call to Replace_Element.	133/2
procedure Replace_Element (Container : in out Vector; Position : in Cursor; New_Item : in Element_Type);	134/2
If Position equals No_Element, then Constraint_Error is propagated; if Position does not designate an element in Container, then Program_Error is propagated. Otherwise Replace_Element assigns New_Item to the element designated by Position. Any exception raised during the assignment is propagated. The element at Position is not an empty element after successful call to Replace_Element.	135/2

```

136/2   procedure Query_Element
        (Container : in Vector;
         Index      : in Index_Type;
         Process    : not null access procedure (Element : in Element_Type));
137/2   If Index is not in the range First_Index (Container) .. Last_Index (Container), then
          Constraint_Error is propagated. Otherwise, Query_Element calls Process.all with the element at
          position Index as the argument. Program_Error is propagated if Process.all tampers with the
          elements of Container. Any exception raised by Process.all is propagated.

138/2   procedure Query_Element
        (Position : in Cursor;
         Process   : not null access procedure (Element : in Element_Type));
139/2   If Position equals No_Element, then Constraint_Error is propagated. Otherwise, Query_Element
          calls Process.all with the element designated by Position as the argument. Program_Error is
          propagated if Process.all tampers with the elements of Container. Any exception raised by
          Process.all is propagated.

140/2   procedure Update_Element
        (Container : in out Vector;
         Index      : in Index_Type;
         Process    : not null access procedure (Element : in out Element_Type));
141/2   If Index is not in the range First_Index (Container) .. Last_Index (Container), then
          Constraint_Error is propagated. Otherwise, Update_Element calls Process.all with the element at
          position Index as the argument. Program_Error is propagated if Process.all tampers with the
          elements of Container. Any exception raised by Process.all is propagated.

142/2   If Element_Type is unconstrained and definite, then the actual Element parameter of Process.all
          shall be unconstrained.

143/2   The element at position Index is not an empty element after successful completion of this
          operation.

144/2   procedure Update_Element
        (Container : in out Vector;
         Position   : in Cursor;
         Process    : not null access procedure (Element : in out Element_Type));
145/2   If Position equals No_Element, then Constraint_Error is propagated; if Position does not
          designate an element in Container, then Program_Error is propagated. Otherwise
          Update_Element calls Process.all with the element designated by Position as the argument.
          Program_Error is propagated if Process.all tampers with the elements of Container. Any
          exception raised by Process.all is propagated.

146/2   If Element_Type is unconstrained and definite, then the actual Element parameter of Process.all
          shall be unconstrained.

147/2   The element designated by Position is not an empty element after successful completion of this
          operation.

148/2   procedure Move (Target : in out Vector;
                    Source  : in out Vector);
149/2   If Target denotes the same object as Source, then Move has no effect. Otherwise, Move first
          calls Clear (Target); then, each element from Source is removed from Source and inserted into
          Target in the original order. The length of Source is 0 after a successful call to Move.

```

```
procedure Insert (Container : in out Vector;
                  Before   : in      Extended_Index;
                  New_Item : in      Vector);
```

150/2

If **Before** is not in the range **First_Index** (**Container**) .. **Last_Index** (**Container**) + 1, then **Constraint_Error** is propagated. If **Length**(**New_Item**) is 0, then **Insert** does nothing. Otherwise, it computes the new length **NL** as the sum of the current length and **Length** (**New_Item**); if the value of **Last** appropriate for length **NL** would be greater than **Index_Type'Last** then **Constraint_Error** is propagated.

151/2

If the current vector capacity is less than **NL**, **Reserve_Capacity** (**Container**, **NL**) is called to increase the vector capacity. Then **Insert** slides the elements in the range **Before** .. **Last_Index** (**Container**) up by **Length**(**New_Item**) positions, and then copies the elements of **New_Item** to the positions starting at **Before**. Any exception raised during the copying is propagated.

152/2

```
procedure Insert (Container : in out Vector;
                  Before   : in      Cursor;
                  New_Item : in      Vector);
```

153/2

If **Before** is not **No_Element**, and does not designate an element in **Container**, then **Program_Error** is propagated. Otherwise, if **Length**(**New_Item**) is 0, then **Insert** does nothing. If **Before** is **No_Element**, then the call is equivalent to **Insert** (**Container**, **Last_Index** (**Container**) + 1, **New_Item**); otherwise the call is equivalent to **Insert** (**Container**, **To_Index** (**Before**), **New_Item**);

154/2

```
procedure Insert (Container : in out Vector;
                  Before   : in      Cursor;
                  New_Item : in      Vector;
                  Position : out    Cursor);
```

155/2

If **Before** is not **No_Element**, and does not designate an element in **Container**, then **Program_Error** is propagated. If **Before** equals **No_Element**, then let **T** be **Last_Index** (**Container**) + 1; otherwise, let **T** be **To_Index** (**Before**). **Insert** (**Container**, **T**, **New_Item**) is called, and then **Position** is set to **To_Cursor** (**Container**, **T**).

156/2

```
procedure Insert (Container : in out Vector;
                  Before   : in      Extended_Index;
                  New_Item : in      Element_Type;
                  Count    : in      Count_Type := 1);
```

157/2

Equivalent to **Insert** (**Container**, **Before**, **To_Vector** (**New_Item**, **Count**));

158/2

```
procedure Insert (Container : in out Vector;
                  Before   : in      Cursor;
                  New_Item : in      Element_Type;
                  Count    : in      Count_Type := 1);
```

159/2

Equivalent to **Insert** (**Container**, **Before**, **To_Vector** (**New_Item**, **Count**));

160/2

```
procedure Insert (Container : in out Vector;
                  Before   : in      Cursor;
                  New_Item : in      Element_Type;
                  Position : out    Cursor;
                  Count    : in      Count_Type := 1);
```

161/2

Equivalent to **Insert** (**Container**, **Before**, **To_Vector** (**New_Item**, **Count**), **Position**);

162/2

- 163/2 **procedure** Insert (Container : **in out** Vector;
 Before : **in** Extended_Index;
 Count : **in** Count_Type := 1);
- 164/2 If Before is not in the range First_Index (Container) .. Last_Index (Container) + 1, then Constraint_Error is propagated. If Count is 0, then Insert does nothing. Otherwise, it computes the new length *NL* as the sum of the current length and Count; if the value of Last appropriate for length *NL* would be greater than Index_Type'Last then Constraint_Error is propagated.
- 165/2 If the current vector capacity is less than *NL*, Reserve_Capacity (Container, *NL*) is called to increase the vector capacity. Then Insert slides the elements in the range Before .. Last_Index (Container) up by Count positions, and then inserts elements that are initialized by default (see 3.3.1) in the positions starting at Before.
- 166/2 **procedure** Insert (Container : **in out** Vector;
 Before : **in** Cursor;
 Position : **out** Cursor;
 Count : **in** Count_Type := 1);
- 167/2 If Before is not No_Element, and does not designate an element in Container, then Program_Error is propagated. If Before equals No_Element, then let *T* be Last_Index (Container) + 1; otherwise, let *T* be To_Index (Before). Insert (Container, *T*, Count) is called, and then Position is set to To_Cursor (Container, *T*).
- 168/2 **procedure** Prepend (Container : **in out** Vector;
 New_Item : **in** Vector;
 Count : **in** Count_Type := 1);
- 169/2 Equivalent to Insert (Container, First_Index (Container), New_Item).
- 170/2 **procedure** Prepend (Container : **in out** Vector;
 New_Item : **in** Element_Type;
 Count : **in** Count_Type := 1);
- 171/2 Equivalent to Insert (Container, First_Index (Container), New_Item, Count).
- 172/2 **procedure** Append (Container : **in out** Vector;
 New_Item : **in** Vector);
- 173/2 Equivalent to Insert (Container, Last_Index (Container) + 1, New_Item).
- 174/2 **procedure** Append (Container : **in out** Vector;
 New_Item : **in** Element_Type;
 Count : **in** Count_Type := 1);
- 175/2 Equivalent to Insert (Container, Last_Index (Container) + 1, New_Item, Count).
- 176/2 **procedure** Insert_Space (Container : **in out** Vector;
 Before : **in** Extended_Index;
 Count : **in** Count_Type := 1);
- 177/2 If Before is not in the range First_Index (Container) .. Last_Index (Container) + 1, then Constraint_Error is propagated. If Count is 0, then Insert_Space does nothing. Otherwise, it computes the new length *NL* as the sum of the current length and Count; if the value of Last appropriate for length *NL* would be greater than Index_Type'Last then Constraint_Error is propagated.
- 178/2 If the current vector capacity is less than *NL*, Reserve_Capacity (Container, *NL*) is called to increase the vector capacity. Then Insert_Space slides the elements in the range Before .. Last_Index (Container) up by Count positions, and then inserts empty elements in the positions starting at Before.

procedure Insert_Space (Container : in out Vector; Before : in Cursor; Position : out Cursor; Count : in Count_Type := 1);	179/2
If Before is not No_Element, and does not designate an element in Container, then Program_Error is propagated. If Before equals No_Element, then let T be Last_Index (Container) + 1; otherwise, let T be To_Index (Before). Insert_Space (Container, T , Count) is called, and then Position is set to To_Cursor (Container, T).	180/2
procedure Delete (Container : in out Vector; Index : in Extended_Index; Count : in Count_Type := 1);	181/2
If Index is not in the range First_Index (Container) .. Last_Index (Container) + 1, then Constraint_Error is propagated. If Count is 0, Delete has no effect. Otherwise Delete slides the elements (if any) starting at position Index + Count down to Index. Any exception raised during element assignment is propagated.	182/2
procedure Delete (Container : in out Vector; Position : in out Cursor; Count : in Count_Type := 1);	183/2
If Position equals No_Element, then Constraint_Error is propagated. If Position does not designate an element in Container, then Program_Error is propagated. Otherwise, Delete (Container, To_Index (Position), Count) is called, and then Position is set to No_Element.	184/2
procedure Delete_First (Container : in out Vector; Count : in Count_Type := 1);	185/2
Equivalent to Delete (Container, First_Index (Container), Count).	186/2
procedure Delete_Last (Container : in out Vector; Count : in Count_Type := 1);	187/2
If Length (Container) <= Count then Delete_Last is equivalent to Clear (Container). Otherwise it is equivalent to Delete (Container, Index_Type'Val(Index_Type'Pos(Last_Index (Container)) – Count + 1), Count).	188/2
procedure Reverse_Elements (Container : in out List);	189/2
Reorders the elements of Container in reverse order.	190/2
procedure Swap (Container : in out Vector; I, J : in Index_Type);	191/2
If either I or J is not in the range First_Index (Container) .. Last_Index (Container), then Constraint_Error is propagated. Otherwise, Swap exchanges the values of the elements at positions I and J.	192/2
procedure Swap (Container : in out Vector; I, J : in Cursor);	193/2
If either I or J is No_Element, then Constraint_Error is propagated. If either I or J do not designate an element in Container, then Program_Error is propagated. Otherwise, Swap exchanges the values of the elements designated by I and J.	194/2
function First_Index (Container : Vector) return Index_Type;	195/2
Returns the value Index_Type'First.	196/2

```

197/2   function First (Container : Vector) return Cursor;
198/2     If Container is empty, First returns No_Element. Otherwise, it returns a cursor that designates
              the first element in Container.

199/2   function First_Element (Container : Vector) return Element_Type;
200/2     Equivalent to Element (Container, First_Index (Container)).

201/2   function Last_Index (Container : Vector) return Extended_Index;
202/2     If Container is empty, Last_Index returns No_Index. Otherwise, it returns the position of the last
              element in Container.

203/2   function Last (Container : Vector) return Cursor;
204/2     If Container is empty, Last returns No_Element. Otherwise, it returns a cursor that designates the
              last element in Container.

205/2   function Last_Element (Container : Vector) return Element_Type;
206/2     Equivalent to Element (Container, Last_Index (Container)).

207/2   function Next (Position : Cursor) return Cursor;
208/2     If Position equals No_Element or designates the last element of the container, then Next returns
              the value No_Element. Otherwise, it returns a cursor that designates the element with index
              To_Index (Position) + 1 in the same vector as Position.

209/2   procedure Next (Position : in out Cursor);
210/2     Equivalent to Position := Next (Position).

211/2   function Previous (Position : Cursor) return Cursor;
212/2     If Position equals No_Element or designates the first element of the container, then Previous
              returns the value No_Element. Otherwise, it returns a cursor that designates the element with
              index To_Index (Position) - 1 in the same vector as Position.

213/2   procedure Previous (Position : in out Cursor);
214/2     Equivalent to Position := Previous (Position).

215/2   function Find_Index (Container : Vector;
                           Item      : Element_Type;
                           Index     : Index_Type := Index_Type'First)
         return Extended_Index;
216/2     Searches the elements of Container for an element equal to Item (using the generic formal
              equality operator). The search starts at position Index and proceeds towards Last_Index
              (Container). If no equal element is found, then Find_Index returns No_Index. Otherwise, it
              returns the index of the first equal element encountered.

217/2   function Find (Container : Vector;
                           Item      : Element_Type;
                           Position  : Cursor := No_Element)
         return Cursor;
218/2     If Position is not No_Element, and does not designate an element in Container, then
              Program_Error is propagated. Otherwise Find searches the elements of Container for an element
              equal to Item (using the generic formal equality operator). The search starts at the first element if
              Position equals No_Element, and at the element designated by Position otherwise. It proceeds

```

towards the last element of Container. If no equal element is found, then Find returns No_Element. Otherwise, it returns a cursor designating the first equal element encountered.

```
function Reverse_Find_Index (Container : Vector;
                             Item      : Element_Type;
                             Index     : Index_Type := Index_Type'Last)
return Extended_Index;
```

219/2

Searches the elements of Container for an element equal to Item (using the generic formal equality operator). The search starts at position Index or, if Index is greater than Last_Index (Container), at position Last_Index (Container). It proceeds towards First_Index (Container). If no equal element is found, then Reverse_Find_Index returns No_Index. Otherwise, it returns the index of the first equal element encountered.

220/2

```
function Reverse_Find (Container : Vector;
                      Item      : Element_Type;
                      Position  : Cursor := No_Element)
return Cursor;
```

221/2

If Position is not No_Element, and does not designate an element in Container, then Program_Error is propagated. Otherwise Reverse_Find searches the elements of Container for an element equal to Item (using the generic formal equality operator). The search starts at the last element if Position equals No_Element, and at the element designated by Position otherwise. It proceeds towards the first element of Container. If no equal element is found, then Reverse_Find returns No_Element. Otherwise, it returns a cursor designating the first equal element encountered.

222/2

```
function Contains (Container : Vector;
                   Item      : Element_Type) return Boolean;
```

223/2

Equivalent to Has_Element (Find (Container, Item)).

224/2

```
function Has_Element (Position : Cursor) return Boolean;
```

225/2

Returns True if Position designates an element, and returns False otherwise.

226/2

```
procedure Iterate
  (Container : in Vector;
   Process    : not null access procedure (Position : in Cursor));
```

227/2

Invokes Process.all with a cursor that designates each element in Container, in index order. Program_Error is propagated if Process.all tampers with the cursors of Container. Any exception raised by Process is propagated.

228/2

```
procedure Reverse_Iterate
  (Container : in Vector;
   Process    : not null access procedure (Position : in Cursor));
```

229/2

Iterates over the elements in Container as per Iterate, except that elements are traversed in reverse index order.

230/2

The actual function for the generic formal function "<" of Generic_Sorting is expected to return the same value each time it is called with a particular pair of element values. It should define a strict ordering relationship, that is, be irreflexive, asymmetric, and transitive; it should not modify Container. If the actual for "<" behaves in some other manner, the behavior of the subprograms of Generic_Sorting are unspecified. How many times the subprograms of Generic_Sorting call "<" is unspecified.

```

232/2   function Is_Sorted (Container : Vector) return Boolean;
233/2     Returns True if the elements are sorted smallest first as determined by the generic formal "<" operator; otherwise, Is_Sorted returns False. Any exception raised during evaluation of "<" is propagated.

234/2   procedure Sort (Container : in out Vector);
235/2     Reorders the elements of Container such that the elements are sorted smallest first as determined by the generic formal "<" operator provided. Any exception raised during evaluation of "<" is propagated.

236/2   procedure Merge (Target  : in out Vector;
237/2           Source   : in out Vector);
238/2     Merge removes elements from Source and inserts them into Target; afterwards, Target contains the union of the elements that were initially in Source and Target; Source is left empty. If Target and Source are initially sorted smallest first, then Target is ordered smallest first as determined by the generic formal "<" operator; otherwise, the order of elements in Target is unspecified. Any exception raised during evaluation of "<" is propagated.

```

Bounded (Run-Time) Errors

- 238/2 Reading the value of an empty element by calling Element, Query_Element, Update_Element, Swap, Is_Sorted, Sort, Merge, "=", Find, or Reverse_Find is a bounded error. The implementation may treat the element as having any normal value (see 13.9.1) of the element type, or raise Constraint_Error or Program_Error before modifying the vector.
- 239/2 Calling Merge in an instance of Generic_Sorting with either Source or Target not ordered smallest first using the provided generic formal "<" operator is a bounded error. Either Program_Error is raised after Target is updated as described for Merge, or the operation works as defined.
- 240/2 A Cursor value is *ambiguous* if any of the following have occurred since it was created:
- 241/2 • Insert, Insert_Space, or Delete has been called on the vector that contains the element the cursor designates with an index value (or a cursor designating an element at such an index value) less than or equal to the index value of the element designated by the cursor; or
 - 242/2 • The vector that contains the element it designates has been passed to the Sort or Merge procedures of an instance of Generic_Sorting, or to the Reverse_Elements procedure.
- 243/2 It is a bounded error to call any subprogram other than "=" or Has_Element declared in Containers.Vectors with an ambiguous (but not invalid, see below) cursor parameter. Possible results are:
- 244/2 • The cursor may be treated as if it were No_Element;
 - 245/2 • The cursor may designate some element in the vector (but not necessarily the element that it originally designated);
 - 246/2 • Constraint_Error may be raised; or
 - 247/2 • Program_Error may be raised.

Erroneous Execution

- 248/2 A Cursor value is *invalid* if any of the following have occurred since it was created:
- 249/2 • The vector that contains the element it designates has been finalized;
 - 250/2 • The vector that contains the element it designates has been used as the Source or Target of a call to Move; or

- The element it designates has been deleted.

251/2

The result of "`=`" or `Has_Element` is unspecified if it is called with an invalid cursor parameter. Execution is erroneous if any other subprogram declared in `Containers.Vectors` is called with an invalid cursor parameter.

Implementation Requirements

No storage associated with a vector object shall be lost upon assignment or scope exit.

253/2

The execution of an `assignment_statement` for a vector shall have the effect of copying the elements from the source vector object to the target vector object.

Implementation Advice

`Containers.Vectors` should be implemented similarly to an array. In particular, if the length of a vector is N , then

- the worst-case time complexity of `Element` should be $O(\log N)$;
- the worst-case time complexity of `Append` with `Count=1` when N is less than the capacity of the vector should be $O(\log N)$; and
- the worst-case time complexity of `Prepend` with `Count=1` and `Delete_First` with `Count=1` should be $O(N \log N)$.

256/2

257

258/2

The worst-case time complexity of a call on procedure `Sort` of an instance of `Containers.Vectors.Generic_Sorting` should be $O(N^{**}2)$, and the average time complexity should be better than $O(N^{**}2)$.

`Containers.Vectors.Generic_Sorting.Sort` and `Containers.Vectors.Generic_Sorting.Merge` should minimize copying of elements.

260/2

`Move` should not copy elements, and should minimize copying of internal data structures.

261/2

If an exception is propagated from a vector operation, no storage should be lost, nor any elements removed from a vector unless specified by the operation.

NOTES

42 All elements of a vector occupy locations in the internal array. If a sparse container is required, a `Hashed_Map` should be used rather than a vector.

263/2

43 If `Index_Type'Base'First = Index_Type'First` an instance of `Ada.Containers.Vectors` will raise `Constraint_Error`. A value below `Index_Type'First` is required so that an empty vector has a meaningful value of `Last_Index`.

264/2

A.18.3 The Package `Containers.Doubly_Linked_Lists`

The language-defined generic package `Containers.Doubly_Linked_Lists` provides private types `List` and `Cursor`, and a set of operations for each type. A list container is optimized for insertion and deletion at any position.

1/2

A doubly-linked list container object manages a linked list of internal *nodes*, each of which contains an element and pointers to the next (successor) and previous (predecessor) internal nodes. A cursor designates a particular node within a list (and by extension the element contained in that node). A cursor keeps designating the same node (and element) as long as the node is part of the container, even if the node is moved in the container.

2/2

The *length* of a list is the number of elements it contains.

3/2

Static Semantics

4/2 The generic library package `Containers.Doubly_Linked_Lists` has the following declaration:

```

5/2   generic
      type Element_Type is private;
      with function "=" (Left, Right : Element_Type)
           return Boolean is <>;
package Ada.Containers.Doubly_Linked_Lists is
  pragma Preelaborate(Doubly_Linked_Lists);

6/2   type List is tagged private;
  pragma Preelaborable_Initialization(List);

7/2   type Cursor is private;
  pragma Preelaborable_Initialization(Cursor);

8/2   Empty_List : constant List;

9/2   No_Element : constant Cursor;

10/2  function "=" (Left, Right : List) return Boolean;

11/2  function Length (Container : List) return Count_Type;

12/2  function Is_Empty (Container : List) return Boolean;

13/2  procedure Clear (Container : in out List);

14/2  function Element (Position : Cursor)
      return Element_Type;

15/2  procedure Replace_Element (Container : in out List;
                                Position  : in      Cursor;
                                New_Item  : in      Element_Type);

16/2  procedure Query_Element
      (Position : in Cursor;
       Process  : not null access procedure (Element : in Element_Type));

17/2  procedure Update_Element
      (Container : in out List;
       Position  : in      Cursor;
       Process   : not null access procedure
                  (Element : in out Element_Type));

18/2  procedure Move (Target : in out List;
                     Source : in out List);

19/2  procedure Insert (Container : in out List;
                      Before   : in      Cursor;
                      New_Item : in      Element_Type;
                      Count    : in      Count_Type := 1);

20/2  procedure Insert (Container : in out List;
                      Before   : in      Cursor;
                      New_Item : in      Element_Type;
                      Position  :        out Cursor;
                      Count    : in      Count_Type := 1);

21/2  procedure Insert (Container : in out List;
                      Before   : in      Cursor;
                      Position  :        out Cursor;
                      Count    : in      Count_Type := 1);

22/2  procedure Prepend (Container : in out List;
                        New_Item : in      Element_Type;
                        Count    : in      Count_Type := 1);

23/2  procedure Append (Container : in out List;
                       New_Item : in      Element_Type;
                       Count    : in      Count_Type := 1);

24/2  procedure Delete (Container : in out List;
                       Position  : in out Cursor;
                       Count    : in      Count_Type := 1);

```

procedure Delete_First (Container : in out List;	25/2
Count : in Count_Type := 1);	
procedure Delete_Last (Container : in out List;	26/2
Count : in Count_Type := 1);	
procedure Reverse_Elements (Container : in out List);	27/2
procedure Swap (Container : in out List;	28/2
I, J : in Cursor);	
procedure Swap_Links (Container : in out List;	29/2
I, J : in Cursor);	
procedure Splice (Target : in out List;	30/2
Before : in Cursor;	
Source : in out List);	
procedure Splice (Target : in out List;	31/2
Before : in Cursor;	
Source : in out List;	
Position : in out Cursor);	
procedure Splice (Container: in out List;	32/2
Before : in Cursor;	
Position : in Cursor);	
function First (Container : List) return Cursor;	33/2
function First_Element (Container : List)	34/2
return Element_Type;	
function Last (Container : List) return Cursor;	35/2
function Last_Element (Container : List)	36/2
return Element_Type;	
function Next (Position : Cursor) return Cursor;	37/2
function Previous (Position : Cursor) return Cursor;	38/2
procedure Next (Position : in out Cursor);	39/2
procedure Previous (Position : in out Cursor);	40/2
function Find (Container : List;	41/2
Item : Element_Type;	
Position : Cursor := No_Element)	
return Cursor;	
function Reverse_Find (Container : List;	42/2
Item : Element_Type;	
Position : Cursor := No_Element)	
return Cursor;	
function Contains (Container : List;	43/2
Item : Element_Type) return Boolean;	
function Has_Element (Position : Cursor) return Boolean;	44/2
procedure Iterate	45/2
(Container : in List;	
Process : not null access procedure (Position : in Cursor));	
procedure Reverse_Iterate	46/2
(Container : in List;	
Process : not null access procedure (Position : in Cursor));	
generic	47/2
with function "<" (Left, Right : Element_Type)	
return Boolean is <>;	
package Generic_Sorting is	
function Is_Sorted (Container : List) return Boolean;	48/2
procedure Sort (Container : in out List);	49/2
procedure Merge (Target : in out List;	50/2
Source : in out List);	

```

51/2      end Generic_Sorting;
52/2  private
53/2    . . . -- not specified by the language
54/2  end Ada.Containers.Doubly_Linked_Lists;
```

55/2 The actual function for the generic formal function "`=`" on `Element_Type` values is expected to define a reflexive and symmetric relationship and return the same result value each time it is called with a particular pair of values. If it behaves in some other manner, the functions `Find`, `Reverse_Find`, and "`=`" on list values return an unspecified value. The exact arguments and number of calls of this generic formal function by the functions `Find`, `Reverse_Find`, and "`=`" on list values are unspecified.

56/2 The type `List` is used to represent lists. The type `List` needs finalization (see 7.6).

57/2 `Empty_List` represents the empty `List` object. It has a length of 0. If an object of type `List` is not otherwise initialized, it is initialized to the same value as `Empty_List`.

58/2 `No_Element` represents a cursor that designates no element. If an object of type `Cursor` is not otherwise initialized, it is initialized to the same value as `No_Element`.

59/2 The predefined "`=`" operator for type `Cursor` returns `True` if both cursors are `No_Element`, or designate the same element in the same container.

60/2 Execution of the default implementation of the `Input`, `Output`, `Read`, or `Write` attribute of type `Cursor` raises `Program_Error`.

61/2 Some operations of this generic package have access-to-subprogram parameters. To ensure such operations are well-defined, they guard against certain actions by the designated subprogram. In particular, some operations check for "tampering with cursors" of a container because they depend on the set of elements of the container remaining constant, and others check for "tampering with elements" of a container because they depend on elements of the container not being replaced.

62/2 A subprogram is said to *tamper with cursors* of a list object `L` if:

- 63/2 • it inserts or deletes elements of `L`, that is, it calls the `Insert`, `Clear`, `Delete`, or `Delete_Last` procedures with `L` as a parameter; or
- 64/2 • it reorders the elements of `L`, that is, it calls the `Splice`, `Swap_Links`, or `Reverse_Elements` procedures or the `Sort` or `Merge` procedures of an instance of `Generic_Sorting` with `L` as a parameter; or
- 65/2 • it finalizes `L`; or
- 66/2 • it calls the `Move` procedure with `L` as a parameter.

67/2 A subprogram is said to *tamper with elements* of a list object `L` if:

- 68/2 • it tampers with cursors of `L`; or
- 69/2 • it replaces one or more elements of `L`, that is, it calls the `Replace_Element` or `Swap` procedures with `L` as a parameter.

70/2 **function** "`=`" (`Left`, `Right` : `List`) **return** Boolean;

71/2 If `Left` and `Right` denote the same list object, then the function returns `True`. If `Left` and `Right` have different lengths, then the function returns `False`. Otherwise, it compares each element in `Left` to the corresponding element in `Right` using the generic formal equality operator. If any such comparison returns `False`, the function returns `False`; otherwise it returns `True`. Any exception raised during evaluation of element equality is propagated.

function Length (Container : List) return Count_Type;	72/2
>Returns the number of elements in Container.	73/2
function Is_Empty (Container : List) return Boolean;	74/2
>Equivalent to Length (Container) = 0.	75/2
procedure Clear (Container : in out List);	76/2
>Removes all the elements from Container.	77/2
function Element (Position : Cursor) return Element_Type;	78/2
>If Position equals No_Element, then Constraint_Error is propagated. Otherwise, Element returns the element designated by Position.	79/2
procedure Replace_Element (Container : in out List; Position : in Cursor; New_Item : in Element_Type);	80/2
>If Position equals No_Element, then Constraint_Error is propagated; if Position does not designate an element in Container, then Program_Error is propagated. Otherwise Replace_Element assigns the value New_Item to the element designated by Position.	81/2
procedure Query_Element (Position : in Cursor; Process : not null access procedure (Element : in Element_Type));	82/2
>If Position equals No_Element, then Constraint_Error is propagated. Otherwise, Query_Element calls Process.all with the element designated by Position as the argument. Program_Error is propagated if Process.all tampers with the elements of Container. Any exception raised by Process.all is propagated.	83/2
procedure Update_Element (Container : in out List; Position : in Cursor; Process : not null access procedure (Element : in out Element_Type));	84/2
>If Position equals No_Element, then Constraint_Error is propagated; if Position does not designate an element in Container, then Program_Error is propagated. Otherwise Update_Element calls Process.all with the element designated by Position as the argument. Program_Error is propagated if Process.all tampers with the elements of Container. Any exception raised by Process.all is propagated.	85/2
>If Element_Type is unconstrained and definite, then the actual Element parameter of Process.all shall be unconstrained.	86/2
procedure Move (Target : in out List; Source : in out List);	87/2
>If Target denotes the same object as Source, then Move has no effect. Otherwise, Move first calls Clear (Target). Then, the nodes in Source are moved to Target (in the original order). The length of Target is set to the length of Source, and the length of Source is set to 0.	88/2
procedure Insert (Container : in out List; Before : in Cursor; New_Item : in Element_Type; Count : in Count_Type := 1);	89/2
>If Before is not No_Element, and does not designate an element in Container, then Program_Error is propagated. Otherwise, Insert inserts Count copies of New_Item prior to the	90/2

element designated by Before. If Before equals No_Element, the new elements are inserted after the last node (if any). Any exception raised during allocation of internal storage is propagated, and Container is not modified.

```
91/2  procedure Insert (Container : in out List;
                      Before    : in      Cursor;
                      New_Item  : in      Element_Type;
                      Position   : out     Cursor;
                      Count     : in      Count_Type := 1);
```

92/2 If Before is not No_Element, and does not designate an element in Container, then Program_Error is propagated. Otherwise, Insert allocates Count copies of New_Item, and inserts them prior to the element designated by Before. If Before equals No_Element, the new elements are inserted after the last element (if any). Position designates the first newly-inserted element. Any exception raised during allocation of internal storage is propagated, and Container is not modified.

```
93/2  procedure Insert (Container : in out List;
                      Before    : in      Cursor;
                      Position  : out     Cursor;
                      Count     : in      Count_Type := 1);
```

94/2 If Before is not No_Element, and does not designate an element in Container, then Program_Error is propagated. Otherwise, Insert inserts Count new elements prior to the element designated by Before. If Before equals No_Element, the new elements are inserted after the last node (if any). The new elements are initialized by default (see 3.3.1). Any exception raised during allocation of internal storage is propagated, and Container is not modified.

```
95/2  procedure Prepend (Container : in out List;
                        New_Item  : in      Element_Type;
                        Count     : in      Count_Type := 1);
```

96/2 Equivalent to Insert (Container, First (Container), New_Item, Count).

```
97/2  procedure Append (Container : in out List;
                        New_Item  : in      Element_Type;
                        Count     : in      Count_Type := 1);
```

98/2 Equivalent to Insert (Container, No_Element, New_Item, Count).

```
99/2  procedure Delete (Container : in out List;
                        Position  : in out Cursor;
                        Count     : in      Count_Type := 1);
```

100/2 If Position equals No_Element, then Constraint_Error is propagated. If Position does not designate an element in Container, then Program_Error is propagated. Otherwise Delete removes (from Container) Count elements starting at the element designated by Position (or all of the elements starting at Position if there are fewer than Count elements starting at Position). Finally, Position is set to No_Element.

```
101/2 procedure Delete_First (Container : in out List;
                            Count      : in      Count_Type := 1);
```

102/2 Equivalent to Delete (Container, First (Container), Count).

```
103/2 procedure Delete_Last (Container : in out List;
                            Count      : in      Count_Type := 1);
```

104/2 If Length (Container) <= Count then Delete_Last is equivalent to Clear (Container). Otherwise it removes the last Count nodes from Container.

procedure Reverse_Elements (Container : in out List);	105/2
Reorders the elements of Container in reverse order.	106/2
procedure Swap (Container : in out List; I, J : in Cursor);	107/2
If either I or J is No_Element, then Constraint_Error is propagated. If either I or J do not designate an element in Container, then Program_Error is propagated. Otherwise, Swap exchanges the values of the elements designated by I and J.	108/2
procedure Swap_Links (Container : in out List; I, J : in Cursor);	109/2
If either I or J is No_Element, then Constraint_Error is propagated. If either I or J do not designate an element in Container, then Program_Error is propagated. Otherwise, Swap_Links exchanges the nodes designated by I and J.	110/2
procedure Splice (Target : in out List; Before : in Cursor; Source : in out List);	111/2
If Before is not No_Element, and does not designate an element in Target, then Program_Error is propagated. Otherwise, if Source denotes the same object as Target, the operation has no effect. Otherwise, Splice reorders elements such that they are removed from Source and moved to Target, immediately prior to Before. If Before equals No_Element, the nodes of Source are spliced after the last node of Target. The length of Target is incremented by the number of nodes in Source, and the length of Source is set to 0.	112/2
procedure Splice (Target : in out List; Before : in Cursor; Source : in out List; Position : in out Cursor);	113/2
If Position is No_Element then Constraint_Error is propagated. If Before does not equal No_Element, and does not designate an element in Target, then Program_Error is propagated. If Position does not equal No_Element, and does not designate a node in Source, then Program_Error is propagated. If Source denotes the same object as Target, then there is no effect if Position equals Before, else the element designated by Position is moved immediately prior to Before, or, if Before equals No_Element, after the last element. In both cases, Position and the length of Target are unchanged. Otherwise the element designated by Position is removed from Source and moved to Target, immediately prior to Before, or, if Before equals No_Element, after the last element of Target. The length of Target is incremented, the length of Source is decremented, and Position is updated to represent an element in Target.	114/2
procedure Splice (Container: in out List; Before : in Cursor; Position : in Cursor);	115/2
If Position is No_Element then Constraint_Error is propagated. If Before does not equal No_Element, and does not designate an element in Container, then Program_Error is propagated. If Position does not equal No_Element, and does not designate a node in Container, then Program_Error is propagated. If Position equals Before there is no effect. Otherwise, the element designated by Position is moved immediately prior to Before, or, if Before equals No_Element, after the last element. The length of Container is unchanged.	116/2

```

117/2   function First (Container : List) return Cursor;
118/2     If Container is empty, First returns the value No_Element. Otherwise it returns a cursor that
             designates the first node in Container.

119/2   function First_Element (Container : List) return Element_Type;
120/2     Equivalent to Element (First (Container)).

121/2   function Last (Container : List) return Cursor;
122/2     If Container is empty, Last returns the value No_Element. Otherwise it returns a cursor that
             designates the last node in Container.

123/2   function Last_Element (Container : List) return Element_Type;
124/2     Equivalent to Element (Last (Container)).

125/2   function Next (Position : Cursor) return Cursor;
126/2     If Position equals No_Element or designates the last element of the container, then Next returns
             the value No_Element. Otherwise, it returns a cursor that designates the successor of the element
             designated by Position.

127/2   function Previous (Position : Cursor) return Cursor;
128/2     If Position equals No_Element or designates the first element of the container, then Previous
             returns the value No_Element. Otherwise, it returns a cursor that designates the predecessor of
             the element designated by Position.

129/2   procedure Next (Position : in out Cursor);
130/2     Equivalent to Position := Next (Position).

131/2   procedure Previous (Position : in out Cursor);
132/2     Equivalent to Position := Previous (Position).

133/2   function Find (Container : List;
                     Item      : Element_Type;
                     Position  : Cursor := No_Element)
             return Cursor;
134/2     If Position is not No_Element, and does not designate an element in Container, then
             Program_Error is propagated. Find searches the elements of Container for an element equal to
             Item (using the generic formal equality operator). The search starts at the element designated by
             Position, or at the first element if Position equals No_Element. It proceeds towards Last
             (Container). If no equal element is found, then Find returns No_Element. Otherwise, it returns a
             cursor designating the first equal element encountered.

135/2   function Reverse_Find (Container : List;
                            Item       : Element_Type;
                            Position   : Cursor := No_Element)
             return Cursor;
136/2     If Position is not No_Element, and does not designate an element in Container, then
             Program_Error is propagated. Find searches the elements of Container for an element equal to
             Item (using the generic formal equality operator). The search starts at the element designated by
             Position, or at the last element if Position equals No_Element. It proceeds towards First
             (Container). If no equal element is found, then Reverse_Find returns No_Element. Otherwise, it
             returns a cursor designating the first equal element encountered.

```

function Contains (Container : List; Item : Element_Type) return Boolean;	137/2
Equivalent to Find (Container, Item) /= No_Element.	138/2
function Has_Element (Position : Cursor) return Boolean;	139/2
Returns True if Position designates an element, and returns False otherwise.	140/2
procedure Iterate (Container : in List; Process : not null access procedure (Position : in Cursor));	141/2
Iterate calls Process.all with a cursor that designates each node in Container, starting with the first node and moving the cursor as per the Next function. Program_Error is propagated if Process.all tampers with the cursors of Container. Any exception raised by Process.all is propagated.	142/2
procedure Reverse_Iterate (Container : in List; Process : not null access procedure (Position : in Cursor));	143/2
Iterates over the nodes in Container as per Iterate, except that elements are traversed in reverse order, starting with the last node and moving the cursor as per the Previous function.	144/2
The actual function for the generic formal function "<" of Generic_Sorting is expected to return the same value each time it is called with a particular pair of element values. It should define a strict ordering relationship, that is, be irreflexive, asymmetric, and transitive; it should not modify Container. If the actual for "<" behaves in some other manner, the behavior of the subprograms of Generic_Sorting are unspecified. How many times the subprograms of Generic_Sorting call "<" is unspecified.	145/2
function Is_Sorted (Container : List) return Boolean;	146/2
Returns True if the elements are sorted smallest first as determined by the generic formal "<" operator; otherwise, Is_Sorted returns False. Any exception raised during evaluation of "<" is propagated.	147/2
procedure Sort (Container : in out List);	148/2
Reorders the nodes of Container such that the elements are sorted smallest first as determined by the generic formal "<" operator provided. The sort is stable. Any exception raised during evaluation of "<" is propagated.	149/2
procedure Merge (Target : in out List; Source : in out List);	150/2
Merge removes elements from Source and inserts them into Target; afterwards, Target contains the union of the elements that were initially in Source and Target; Source is left empty. If Target and Source are initially sorted smallest first, then Target is ordered smallest first as determined by the generic formal "<" operator; otherwise, the order of elements in Target is unspecified. Any exception raised during evaluation of "<" is propagated.	151/2

Bounded (Run-Time) Errors

Calling Merge in an instance of Generic_Sorting with either Source or Target not ordered smallest first using the provided generic formal "<" operator is a bounded error. Either Program_Error is raised after Target is updated as described for Merge, or the operation works as defined.

Erroneous Execution

- 153/2 A Cursor value is *invalid* if any of the following have occurred since it was created:
- The list that contains the element it designates has been finalized;
 - The list that contains the element it designates has been used as the Source or Target of a call to Move; or
 - The element it designates has been deleted.
- 157/2 The result of " $=$ " or Has_Element is unspecified if it is called with an invalid cursor parameter. Execution is erroneous if any other subprogram declared in Containers.Doubly_Linked_Lists is called with an invalid cursor parameter.

Implementation Requirements

- 158/2 No storage associated with a doubly-linked List object shall be lost upon assignment or scope exit.
- 159/2 The execution of an **assignment_statement** for a list shall have the effect of copying the elements from the source list object to the target list object.

Implementation Advice

- 160/2 Containers.Doubly_Linked_Lists should be implemented similarly to a linked list. In particular, if N is the length of a list, then the worst-case time complexity of Element, Insert with Count=1, and Delete with Count=1 should be $O(\log N)$.
- 161/2 The worst-case time complexity of a call on procedure Sort of an instance of Containers.Doubly_Linked_Lists.Generic_Sorting should be $O(N^{**2})$, and the average time complexity should be better than $O(N^{**2})$.
- 162/2 Move should not copy elements, and should minimize copying of internal data structures.
- 163/2 If an exception is propagated from a list operation, no storage should be lost, nor any elements removed from a list unless specified by the operation.

NOTES

- 164/2 44 Sorting a list never copies elements, and is a stable sort (equal elements remain in the original order). This is different than sorting an array or vector, which may need to copy elements, and is probably not a stable sort.

A.18.4 Maps

- 1/2 The language-defined generic packages Containers.Hashed_Maps and Containers.Ordered_Maps provide private types Map and Cursor, and a set of operations for each type. A map container allows an arbitrary type to be used as a key to find the element associated with that key. A hashed map uses a hash function to organize the keys, while an ordered map orders the keys per a specified relation.
- 2/2 This section describes the declarations that are common to both kinds of maps. See A.18.5 for a description of the semantics specific to Containers.Hashed_Maps and A.18.6 for a description of the semantics specific to Containers.Ordered_Maps.

Static Semantics

- 3/2 The actual function for the generic formal function " $=$ " on Element_Type values is expected to define a reflexive and symmetric relationship and return the same result value each time it is called with a particular pair of values. If it behaves in some other manner, the function " $=$ " on map values returns an

unspecified value. The exact arguments and number of calls of this generic formal function by the function "`=`" on map values are unspecified.

The type Map is used to represent maps. The type Map needs finalization (see 7.6). 4/2

A map contains pairs of keys and elements, called *nodes*. Map cursors designate nodes, but also can be thought of as designating an element (the element contained in the node) for consistency with the other containers. There exists an equivalence relation on keys, whose definition is different for hashed maps and ordered maps. A map never contains two or more nodes with equivalent keys. The *length* of a map is the number of nodes it contains. 5/2

Each nonempty map has two particular nodes called the *first node* and the *last node* (which may be the same). Each node except for the last node has a *successor node*. If there are no other intervening operations, starting with the first node and repeatedly going to the successor node will visit each node in the map exactly once until the last node is reached. The exact definition of these terms is different for hashed maps and ordered maps. 6/2

Some operations of these generic packages have access-to-subprogram parameters. To ensure such operations are well-defined, they guard against certain actions by the designated subprogram. In particular, some operations check for "tampering with cursors" of a container because they depend on the set of elements of the container remaining constant, and others check for "tampering with elements" of a container because they depend on elements of the container not being replaced. 7/2

A subprogram is said to *tamper with cursors* of a map object *M* if: 8/2

- it inserts or deletes elements of *M*, that is, it calls the Insert, Include, Clear, Delete, or Exclude procedures with *M* as a parameter; or 9/2
- it finalizes *M*; or 10/2
- it calls the Move procedure with *M* as a parameter; or 11/2
- it calls one of the operations defined to tamper with the cursors of *M*. 12/2

A subprogram is said to *tamper with elements* of a map object *M* if: 13/2

- it tampers with cursors of *M*; or 14/2
- it replaces one or more elements of *M*, that is, it calls the Replace or Replace_Element procedures with *M* as a parameter. 15/2

`Empty_Map` represents the empty Map object. It has a length of 0. If an object of type Map is not otherwise initialized, it is initialized to the same value as `Empty_Map`. 16/2

`No_Element` represents a cursor that designates no node. If an object of type Cursor is not otherwise initialized, it is initialized to the same value as `No_Element`. 17/2

The predefined "`=`" operator for type Cursor returns True if both cursors are `No_Element`, or designate the same element in the same container. 18/2

Execution of the default implementation of the Input, Output, Read, or Write attribute of type Cursor raises Program_Error. 19/2

```
function "=" (Left, Right : Map) return Boolean;
```

If Left and Right denote the same map object, then the function returns True. If Left and Right have different lengths, then the function returns False. Otherwise, for each key *K* in Left, the function returns False if:

- a key equivalent to *K* is not present in Right; or 22/2

- 23/2 • the element associated with K in Left is not equal to the element associated with K in Right (using the generic formal equality operator for elements).
- 24/2 If the function has not returned a result after checking all of the keys, it returns True. Any exception raised during evaluation of key equivalence or element equality is propagated.
- 25/2 **function** Length (Container : Map) **return** Count_Type;
- 26/2 Returns the number of nodes in Container.
- 27/2 **function** Is_Empty (Container : Map) **return** Boolean;
- 28/2 Equivalent to Length (Container) = 0.
- 29/2 **procedure** Clear (Container : **in out** Map);
- 30/2 Removes all the nodes from Container.
- 31/2 **function** Key (Position : Cursor) **return** Key_Type;
- 32/2 If Position equals No_Element, then Constraint_Error is propagated. Otherwise, Key returns the key component of the node designated by Position.
- 33/2 **function** Element (Position : Cursor) **return** Element_Type;
- 34/2 If Position equals No_Element, then Constraint_Error is propagated. Otherwise, Element returns the element component of the node designated by Position.
- 35/2 **procedure** Replace_Element (Container : **in out** Map;
 Position : **in** Cursor;
 New_Item : **in** Element_Type);
- 36/2 If Position equals No_Element, then Constraint_Error is propagated; if Position does not designate an element in Container, then Program_Error is propagated. Otherwise Replace_Element assigns New_Item to the element of the node designated by Position.
- 37/2 **procedure** Query_Element
 (Position : **in** Cursor;
 Process : **not null access procedure** (Key : **in** Key_Type;
 Element : **in** Element_Type));
- 38/2 If Position equals No_Element, then Constraint_Error is propagated. Otherwise, Query_Element calls Process.all with the key and element from the node designated by Position as the arguments. Program_Error is propagated if Process.all tampers with the elements of Container. Any exception raised by Process.all is propagated.
- 39/2 **procedure** Update_Element
 (Container : **in out** Map;
 Position : **in** Cursor;
 Process : **not null access procedure** (Key : **in** Key_Type;
 Element : **in out** Element_Type));
- 40/2 If Position equals No_Element, then Constraint_Error is propagated; if Position does not designate an element in Container, then Program_Error is propagated. Otherwise Update_Element calls Process.all with the key and element from the node designated by Position as the arguments. Program_Error is propagated if Process.all tampers with the elements of Container. Any exception raised by Process.all is propagated.
- 41/2 If Element_Type is unconstrained and definite, then the actual Element parameter of Process.all shall be unconstrained.

```
procedure Move (Target : in out Map;
                 Source : in out Map);
```

42/2

If Target denotes the same object as Source, then Move has no effect. Otherwise, Move first calls Clear (Target). Then, each node from Source is removed from Source and inserted into Target. The length of Source is 0 after a successful call to Move.

```
procedure Insert (Container : in out Map;
                  Key      : in      Key_Type;
                  New_Item : in      Element_Type;
                  Position : out    Cursor;
                  Inserted : out    Boolean);
```

44/2

Insert checks if a node with a key equivalent to Key is already present in Container. If a match is found, Inserted is set to False and Position designates the element with the matching key. Otherwise, Insert allocates a new node, initializes it to Key and New_Item, and adds it to Container; Inserted is set to True and Position designates the newly-inserted node. Any exception raised during allocation is propagated and Container is not modified.

45/2

```
procedure Insert (Container : in out Map;
                  Key      : in      Key_Type;
                  Position : out    Cursor;
                  Inserted : out    Boolean);
```

46/2

Insert inserts Key into Container as per the five-parameter Insert, with the difference that an element initialized by default (see 3.3.1) is inserted.

47/2

```
procedure Insert (Container : in out Map;
                  Key      : in      Key_Type;
                  New_Item : in      Element_Type);
```

48/2

Insert inserts Key and New_Item into Container as per the five-parameter Insert, with the difference that if a node with a key equivalent to Key is already in the map, then Constraint_Error is propagated.

49/2

```
procedure Include (Container : in out Map;
                   Key      : in      Key_Type;
                   New_Item : in      Element_Type);
```

50/2

Include inserts Key and New_Item into Container as per the five-parameter Insert, with the difference that if a node with a key equivalent to Key is already in the map, then this operation assigns Key and New_Item to the matching node. Any exception raised during assignment is propagated.

51/2

```
procedure Replace (Container : in out Map;
                   Key      : in      Key_Type;
                   New_Item : in      Element_Type);
```

52/2

Replace checks if a node with a key equivalent to Key is present in Container. If a match is found, Replace assigns Key and New_Item to the matching node; otherwise, Constraint_Error is propagated.

53/2

```
procedure Exclude (Container : in out Map;
                   Key      : in      Key_Type);
```

54/2

Exclude checks if a node with a key equivalent to Key is present in Container. If a match is found, Exclude removes the node from the map.

55/2

```

56/2   procedure Delete (Container : in out Map;
                      Key      : in      Key_Type);
57/2     Delete checks if a node with a key equivalent to Key is present in Container. If a match is found,
             Delete removes the node from the map; otherwise, Constraint_Error is propagated.

58/2   procedure Delete (Container : in out Map;
                      Position  : in out Cursor);
59/2     If Position equals No_Element, then Constraint_Error is propagated. If Position does not
             designate an element in Container, then Program_Error is propagated. Otherwise, Delete
             removes the node designated by Position from the map. Position is set to No_Element on return.

60/2   function First (Container : Map) return Cursor;
61/2     If Length (Container) = 0, then First returns No_Element. Otherwise, First returns a cursor that
             designates the first node in Container.

62/2   function Next (Position : Cursor) return Cursor;
63/2     Returns a cursor that designates the successor of the node designated by Position. If Position
             designates the last node, then No_Element is returned. If Position equals No_Element, then
             No_Element is returned.

64/2   procedure Next (Position : in out Cursor);
65/2     Equivalent to Position := Next (Position).

66/2   function Find (Container : Map;
                      Key      : Key_Type) return Cursor;
67/2     If Length (Container) equals 0, then Find returns No_Element. Otherwise, Find checks if a node
             with a key equivalent to Key is present in Container. If a match is found, a cursor designating
             the matching node is returned; otherwise, No_Element is returned.

68/2   function Element (Container : Map;
                      Key      : Key_Type) return Element_Type;
69/2     Equivalent to Element (Find (Container, Key)).

70/2   function Contains (Container : Map;
                      Key      : Key_Type) return Boolean;
71/2     Equivalent to Find (Container, Key) /= No_Element.

72/2   function Has_Element (Position : Cursor) return Boolean;
73/2     Returns True if Position designates a node, and returns False otherwise.

74/2   procedure Iterate
          (Container : in Map;
           Process   : not null access procedure (Position : in Cursor));
75/2     Iterate calls Process.all with a cursor that designates each node in Container, starting with the
             first node and moving the cursor according to the successor relation. Program_Error is
             propagated if Process.all tampers with the cursors of Container. Any exception raised by
             Process.all is propagated.

```

Erroneous Execution

76/2 A Cursor value is *invalid* if any of the following have occurred since it was created:

- The map that contains the node it designates has been finalized;

- The map that contains the node it designates has been used as the Source or Target of a call to Move; or
- The node it designates has been deleted from the map.

78/2

79/2

The result of "`=`" or `Has_Element` is unspecified if these functions are called with an invalid cursor parameter. Execution is erroneous if any other subprogram declared in `Containers.Hashed_Maps` or `Containers.Ordered_Maps` is called with an invalid cursor parameter.

80/2

Implementation Requirements

No storage associated with a Map object shall be lost upon assignment or scope exit.

81/2

The execution of an `assignment_statement` for a map shall have the effect of copying the elements from the source map object to the target map object.

82/2

Implementation Advice

Move should not copy elements, and should minimize copying of internal data structures.

83/2

If an exception is propagated from a map operation, no storage should be lost, nor any elements removed from a map unless specified by the operation.

84/2

A.18.5 The Package `Containers.Hashed_Maps`

Static Semantics

The generic library package `Containers.Hashed_Maps` has the following declaration:

1/2

```
generic
  type Key_Type is private;
  type Element_Type is private;
  with function Hash (Key : Key_Type) return Hash_Type;
  with function Equivalent_Keys (Left, Right : Key_Type)
    return Boolean;
  with function "=" (Left, Right : Element_Type)
    return Boolean is <>;
package Ada.Containers.Hashed_Maps is
  pragma Preelaborate(Hashed_Maps);
  type Map is tagged private;
  pragma Preelaborate_Initialization(Map);
  type Cursor is private;
  pragma Preelaborate_Initialization(Cursor);
  Empty_Map : constant Map;
  No_Element : constant Cursor;
  function "=" (Left, Right : Map) return Boolean;
  function Capacity (Container : Map) return Count_Type;
  procedure Reserve_Capacity (Container : in out Map;
                               Capacity : in      Count_Type);
  function Length (Container : Map) return Count_Type;
  function Is_Empty (Container : Map) return Boolean;
  procedure Clear (Container : in out Map);
  function Key (Position : Cursor) return Key_Type;
  function Element (Position : Cursor) return Element_Type;
  procedure Replace_Element (Container : in out Map;
                            Position   : in      Cursor;
                            New_Item   : in      Element_Type);

```

2/2

3/2

4/2

5/2

6/2

7/2

8/2

9/2

10/2

11/2

12/2

13/2

14/2

15/2

```

16/2      procedure Query_Element
          (Position : in Cursor;
           Process  : not null access procedure (Key      : in Key_Type;
                                                    Element : in Element_Type));
17/2      procedure Update_Element
          (Container : in out Map;
           Position  : in      Cursor;
           Process   : not null access procedure
                      (Key      : in      Key_Type;
                       Element : in out Element_Type));
18/2      procedure Move (Target : in out Map;
                      Source : in out Map);
19/2      procedure Insert (Container : in out Map;
                        Key       : in      Key_Type;
                        New_Item : in      Element_Type;
                        Position  : out    Cursor;
                        Inserted  : out    Boolean);
20/2      procedure Insert (Container : in out Map;
                        Key       : in      Key_Type;
                        Position  : out    Cursor;
                        Inserted  : out    Boolean);
21/2      procedure Insert (Container : in out Map;
                        Key       : in      Key_Type;
                        New_Item : in      Element_Type);
22/2      procedure Include (Container : in out Map;
                           Key       : in      Key_Type;
                           New_Item : in      Element_Type);
23/2      procedure Replace (Container : in out Map;
                           Key       : in      Key_Type;
                           New_Item : in      Element_Type);
24/2      procedure Exclude (Container : in out Map;
                           Key       : in      Key_Type);
25/2      procedure Delete (Container : in out Map;
                           Key       : in      Key_Type);
26/2      procedure Delete (Container : in out Map;
                           Position  : in out Cursor);
27/2      function First (Container : Map)
                  return Cursor;
28/2      function Next (Position  : Cursor) return Cursor;
29/2      procedure Next (Position  : in out Cursor);
30/2      function Find (Container : Map;
                        Key       : Key_Type)
                  return Cursor;
31/2      function Element (Container : Map;
                           Key       : Key_Type)
                  return Element_Type;
32/2      function Contains (Container : Map;
                           Key       : Key_Type) return Boolean;
33/2      function Has_Element (Position : Cursor) return Boolean;
34/2      function Equivalent_Keys (Left, Right : Cursor)
                  return Boolean;
35/2      function Equivalent_Keys (Left   : Cursor;
                                 Right  : Key_Type)
                  return Boolean;
36/2      function Equivalent_Keys (Left   : Key_Type;
                                 Right  : Cursor)
                  return Boolean;

```

```

procedure Iterate
  (Container : in Map;
   Process   : not null access procedure (Position : in Cursor));
private
  ... -- not specified by the language
end Ada.Containers.Hashed_Maps;

```

An object of type Map contains an expandable hash table, which is used to provide direct access to nodes. The *capacity* of an object of type Map is the maximum number of nodes that can be inserted into the hash table prior to it being automatically expanded.

Two keys K_1 and K_2 are defined to be *equivalent* if Equivalent_Keys(K_1, K_2) returns True. 42/2

The actual function for the generic formal function Hash is expected to return the same value each time it is called with a particular key value. For any two equivalent key values, the actual for Hash is expected to return the same value. If the actual for Hash behaves in some other manner, the behavior of this package is unspecified. Which subprograms of this package call Hash, and how many times they call it, is unspecified. 43/2

The actual function for the generic formal function Equivalent_Keys on Key_Type values is expected to return the same value each time it is called with a particular pair of key values. It should define an equivalence relationship, that is, be reflexive, symmetric, and transitive. If the actual for Equivalent_Keys behaves in some other manner, the behavior of this package is unspecified. Which subprograms of this package call Equivalent_Keys, and how many times they call it, is unspecified. 44/2

If the value of a key stored in a node of a map is changed other than by an operation in this package such that at least one of Hash or Equivalent_Keys give different results, the behavior of this package is unspecified. 45/2

Which nodes are the first node and the last node of a map, and which node is the successor of a given node, are unspecified, other than the general semantics described in A.18.4. 46/2

```

function Capacity (Container : Map) return Count_Type; 47/2
  Returns the capacity of Container. 48/2

```

```

procedure Reserve_Capacity (Container : in out Map;
                            Capacity   : in      Count_Type); 49/2

```

Reserve_Capacity allocates a new hash table such that the length of the resulting map can become at least the value Capacity without requiring an additional call to Reserve_Capacity, and is large enough to hold the current length of Container. Reserve_Capacity then rehashes the nodes in Container onto the new hash table. It replaces the old hash table with the new hash table, and then deallocates the old hash table. Any exception raised during allocation is propagated and Container is not modified. 50/2

Reserve_Capacity tampers with the cursors of Container. 51/2

```

procedure Clear (Container : in out Map); 52/2

```

In addition to the semantics described in A.18.4, Clear does not affect the capacity of Container. 53/2

```

54/2      procedure Insert (Container : in out Map;
                           Key      : in      Key_Type;
                           New_Item : in      Element_Type;
                           Position  :        out Cursor;
                           Inserted   :        out Boolean);

55/2      In addition to the semantics described in A.18.4, if Length (Container) equals Capacity
           (Container), then Insert first calls Reserve_Capacity to increase the capacity of Container to
           some larger value.

56/2      function Equivalent_Keys (Left, Right : Cursor)
                           return Boolean;

57/2      Equivalent to Equivalent_Keys (Key (Left), Key (Right)).

58/2      function Equivalent_Keys (Left : Cursor;
                           Right : Key_Type) return Boolean;

59/2      Equivalent to Equivalent_Keys (Key (Left), Right).

60/2      function Equivalent_Keys (Left : Key_Type;
                           Right : Cursor) return Boolean;

61/2      Equivalent to Equivalent_Keys (Left, Key (Right)).

```

Implementation Advice

62/2 If N is the length of a map, the average time complexity of the subprograms Element, Insert, Include, Replace, Delete, Exclude and Find that take a key parameter should be $O(\log N)$. The average time complexity of the subprograms that take a cursor parameter should be $O(1)$. The average time complexity of Reserve_Capacity should be $O(N)$.

A.18.6 The Package Containers.Ordered_Maps

Static Semantics

1/2 The generic library package Containers.Ordered_Maps has the following declaration:

```

2/2      generic
              type Key_Type is private;
              type Element_Type is private;
              with function "<" (Left, Right : Key_Type) return Boolean is <>;
              with function "=" (Left, Right : Element_Type) return Boolean is <>;
      package Ada.Containers.Ordered_Maps is
          pragma Preelaborate(Ordered_Maps);

3/2          function Equivalent_Keys (Left, Right : Key_Type) return Boolean;

4/2          type Map is tagged private;
          pragma Preelaborable_Initialization(Map);

5/2          type Cursor is private;
          pragma Preelaborable_Initialization(Cursor);

6/2          Empty_Map : constant Map;

7/2          No_Element : constant Cursor;

8/2          function "=" (Left, Right : Map) return Boolean;

9/2          function Length (Container : Map) return Count_Type;

10/2         function Is_Empty (Container : Map) return Boolean;

11/2        procedure Clear (Container : in out Map);

12/2        function Key (Position : Cursor) return Key_Type;

13/2        function Element (Position : Cursor) return Element_Type;

```

procedure Replace_Element (Container : in out Map; Position : in Cursor; New_Item : in Element_Type);	14/2
procedure Query_Element (Position : in Cursor; Process : not null access procedure (Key : in Key_Type; Element : in out Element_Type));	15/2
procedure Update_Element (Container : in out Map; Position : in Cursor; Process : not null access procedure (Key : in Key_Type; Element : in out Element_Type));	16/2
procedure Move (Target : in out Map; Source : in out Map);	17/2
procedure Insert (Container : in out Map; Key : in Key_Type; New_Item : in Element_Type; Position : out Cursor; Inserted : out Boolean);	18/2
procedure Insert (Container : in out Map; Key : in Key_Type; Position : out Cursor; Inserted : out Boolean);	19/2
procedure Insert (Container : in out Map; Key : in Key_Type; New_Item : in Element_Type);	20/2
procedure Include (Container : in out Map; Key : in Key_Type; New_Item : in Element_Type);	21/2
procedure Replace (Container : in out Map; Key : in Key_Type; New_Item : in Element_Type);	22/2
procedure Exclude (Container : in out Map; Key : in Key_Type);	23/2
procedure Delete (Container : in out Map; Key : in Key_Type);	24/2
procedure Delete (Container : in out Map; Position : in out Cursor);	25/2
procedure Delete_First (Container : in out Map);	26/2
procedure Delete_Last (Container : in out Map);	27/2
function First (Container : Map) return Cursor;	28/2
function First_Element (Container : Map) return Element_Type;	29/2
function First_Key (Container : Map) return Key_Type;	30/2
function Last (Container : Map) return Cursor;	31/2
function Last_Element (Container : Map) return Element_Type;	32/2
function Last_Key (Container : Map) return Key_Type;	33/2
function Next (Position : Cursor) return Cursor;	34/2
procedure Next (Position : in out Cursor);	35/2
function Previous (Position : Cursor) return Cursor;	36/2
procedure Previous (Position : in out Cursor);	37/2
function Find (Container : Map; Key : Key_Type) return Cursor;	38/2
function Element (Container : Map; Key : Key_Type) return Element_Type;	39/2

```

40/2      function Floor (Container : Map;
                      Key       : Key_Type) return Cursor;
41/2      function Ceiling (Container : Map;
                           Key       : Key_Type) return Cursor;
42/2      function Contains (Container : Map;
                           Key       : Key_Type) return Boolean;
43/2      function Has_Element (Position : Cursor) return Boolean;
44/2      function "<" (Left, Right : Cursor) return Boolean;
45/2      function ">" (Left, Right : Cursor) return Boolean;
46/2      function "<" (Left : Cursor; Right : Key_Type) return Boolean;
47/2      function ">" (Left : Cursor; Right : Key_Type) return Boolean;
48/2      function "<" (Left : Key_Type; Right : Cursor) return Boolean;
49/2      function ">" (Left : Key_Type; Right : Cursor) return Boolean;
50/2      procedure Iterate
              (Container : in Map;
               Process   : not null access procedure (Position : in Cursor));
51/2      procedure Reverse_Iterate
              (Container : in Map;
               Process   : not null access procedure (Position : in Cursor));
52/2      private
53/2          ... -- not specified by the language
54/2      end Ada.Containers.Ordered_Maps;

55/2 Two keys  $K_1$  and  $K_2$  are equivalent if both  $K_1 < K_2$  and  $K_2 < K_1$  return False, using the generic formal " $<$ " operator for keys. Function Equivalent_Keys returns True if Left and Right are equivalent, and False otherwise.

56/2 The actual function for the generic formal function " $<$ " on Key_Type values is expected to return the same value each time it is called with a particular pair of key values. It should define a strict ordering relationship, that is, be irreflexive, asymmetric, and transitive. If the actual for " $<$ " behaves in some other manner, the behavior of this package is unspecified. Which subprograms of this package call " $<$ " and how many times they call it, is unspecified.

57/2 If the value of a key stored in a map is changed other than by an operation in this package such that at least one of " $<$ " or " $=$ " give different results, the behavior of this package is unspecified.

58/2 The first node of a nonempty map is the one whose key is less than the key of all the other nodes in the map. The last node of a nonempty map is the one whose key is greater than the key of all the other elements in the map. The successor of a node is the node with the smallest key that is larger than the key of the given node. The predecessor of a node is the node with the largest key that is smaller than the key of the given node. All comparisons are done using the generic formal " $<$ " operator for keys.

59/2      procedure Delete_First (Container : in out Map);
60/2          If Container is empty, Delete_First has no effect. Otherwise the node designated by First (Container) is removed from Container. Delete_First tampers with the cursors of Container.

61/2      procedure Delete_Last (Container : in out Map);
62/2          If Container is empty, Delete_Last has no effect. Otherwise the node designated by Last (Container) is removed from Container. Delete_Last tampers with the cursors of Container.

63/2      function First_Element (Container : Map) return Element_Type;
64/2          Equivalent to Element (First (Container)).

```

function First_Key (Container : Map) return Key_Type;	65/2
Equivalent to Key (First (Container)).	66/2
function Last (Container : Map) return Cursor;	67/2
Returns a cursor that designates the last node in Container. If Container is empty, returns No_Element.	68/2
function Last_Element (Container : Map) return Element_Type;	69/2
Equivalent to Element (Last (Container)).	70/2
function Last_Key (Container : Map) return Key_Type;	71/2
Equivalent to Key (Last (Container)).	72/2
function Previous (Position : Cursor) return Cursor;	73/2
If Position equals No_Element, then Previous returns No_Element. Otherwise Previous returns a cursor designating the node that precedes the one designated by Position. If Position designates the first element, then Previous returns No_Element.	74/2
procedure Previous (Position : in out Cursor);	75/2
Equivalent to Position := Previous (Position).	76/2
function Floor (Container : Map; Key : Key_Type) return Cursor;	77/2
Floor searches for the last node whose key is not greater than Key, using the generic formal "<" operator for keys. If such a node is found, a cursor that designates it is returned. Otherwise No_Element is returned.	78/2
function Ceiling (Container : Map; Key : Key_Type) return Cursor;	79/2
Ceiling searches for the first node whose key is not less than Key, using the generic formal "<" operator for keys. If such a node is found, a cursor that designates it is returned. Otherwise No_Element is returned.	80/2
function "<" (Left, Right : Cursor) return Boolean;	81/2
Equivalent to Key (Left) < Key (Right).	82/2
function ">" (Left, Right : Cursor) return Boolean;	83/2
Equivalent to Key (Right) < Key (Left).	84/2
function "<" (Left : Cursor; Right : Key_Type) return Boolean;	85/2
Equivalent to Key (Left) < Right.	86/2
function ">" (Left : Cursor; Right : Key_Type) return Boolean;	87/2
Equivalent to Right < Key (Left).	88/2
function "<" (Left : Key_Type; Right : Cursor) return Boolean;	89/2
Equivalent to Left < Key (Right).	90/2
function ">" (Left : Key_Type; Right : Cursor) return Boolean;	91/2
Equivalent to Key (Right) < Left.	92/2

93/2 **procedure** Reverse_Iterate
 (Container : **in** Map;
 Process : **not null access procedure** (Position : **in** Cursor));
 94/2 Iterates over the nodes in Container as per Iterate, with the difference that the nodes are traversed in predecessor order, starting with the last node.

Implementation Advice

95/2 If N is the length of a map, then the worst-case time complexity of the Element, Insert, Include, Replace, Delete, Exclude and Find operations that take a key parameter should be $O((\log N)^{**}2)$ or better. The worst-case time complexity of the subprograms that take a cursor parameter should be $O(1)$.

A.18.7 Sets

- 1/2 The language-defined generic packages Containers.Hashed_Sets and Containers.Ordered_Sets provide private types Set and Cursor, and a set of operations for each type. A set container allows elements of an arbitrary type to be stored without duplication. A hashed set uses a hash function to organize elements, while an ordered set orders its element per a specified relation.
- 2/2 This section describes the declarations that are common to both kinds of sets. See A.18.8 for a description of the semantics specific to Containers.Hashed_Sets and A.18.9 for a description of the semantics specific to Containers.Ordered_Sets.

Static Semantics

- 3/2 The actual function for the generic formal function " $=$ " on Element_Type values is expected to define a reflexive and symmetric relationship and return the same result value each time it is called with a particular pair of values. If it behaves in some other manner, the function " $=$ " on set values returns an unspecified value. The exact arguments and number of calls of this generic formal function by the function " $=$ " on set values are unspecified.
- 4/2 The type Set is used to represent sets. The type Set needs finalization (see 7.6).
- 5/2 A set contains elements. Set cursors designate elements. There exists an equivalence relation on elements, whose definition is different for hashed sets and ordered sets. A set never contains two or more equivalent elements. The *length* of a set is the number of elements it contains.
- 6/2 Each nonempty set has two particular elements called the *first element* and the *last element* (which may be the same). Each element except for the last element has a *successor element*. If there are no other intervening operations, starting with the first element and repeatedly going to the successor element will visit each element in the set exactly once until the last element is reached. The exact definition of these terms is different for hashed sets and ordered sets.
- 7/2 Some operations of these generic packages have access-to-subprogram parameters. To ensure such operations are well-defined, they guard against certain actions by the designated subprogram. In particular, some operations check for "tampering with cursors" of a container because they depend on the set of elements of the container remaining constant, and others check for "tampering with elements" of a container because they depend on elements of the container not being replaced.
- 8/2 A subprogram is said to *tamper with cursors* of a set object S if:
- 9/2
 - it inserts or deletes elements of S , that is, it calls the Insert, Include, Clear, Delete, Exclude, or Replace_Element procedures with S as a parameter; or
 - it finalizes S ; or

- it calls the Move procedure with S as a parameter; or 11/2
- it calls one of the operations defined to tamper with cursors of S . 12/2

A subprogram is said to *tamper with elements* of a set object S if: 13/2

- it tampers with cursors of S . 14/2

Empty_Set represents the empty Set object. It has a length of 0. If an object of type Set is not otherwise initialized, it is initialized to the same value as Empty_Set. 15/2

No_Element represents a cursor that designates no element. If an object of type Cursor is not otherwise initialized, it is initialized to the same value as No_Element. 16/2

The predefined " $=$ " operator for type Cursor returns True if both cursors are No_Element, or designate the same element in the same container. 17/2

Execution of the default implementation of the Input, Output, Read, or Write attribute of type Cursor raises Program_Error. 18/2

```
function " $=$ " (Left, Right : Set) return Boolean;
```

If Left and Right denote the same set object, then the function returns True. If Left and Right have different lengths, then the function returns False. Otherwise, for each element E in Left, the function returns False if an element equal to E (using the generic formal equality operator) is not present in Right. If the function has not returned a result after checking all of the elements, it returns True. Any exception raised during evaluation of element equality is propagated. 20/2

```
function Equivalent_Sets (Left, Right : Set) return Boolean;
```

If Left and Right denote the same set object, then the function returns True. If Left and Right have different lengths, then the function returns False. Otherwise, for each element E in Left, the function returns False if an element equivalent to E is not present in Right. If the function has not returned a result after checking all of the elements, it returns True. Any exception raised during evaluation of element equivalence is propagated. 22/2

```
function To_Set (New_Item : Element_Type) return Set;
```

Returns a set containing the single element New_Item. 24/2

```
function Length (Container : Set) return Count_Type;
```

Returns the number of elements in Container. 26/2

```
function Is_Empty (Container : Set) return Boolean;
```

Equivalent to Length (Container) = 0. 28/2

```
procedure Clear (Container : in out Set);
```

Removes all the elements from Container. 30/2

```
function Element (Position : Cursor) return Element_Type;
```

If Position equals No_Element, then Constraint_Error is propagated. Otherwise, Element returns the element designated by Position. 32/2

33/2 **procedure** Replace_Element (*Container* : **in out** Set;
 Position : **in** Cursor;
 New_Item : **in** Element_Type);

34/2 If *Position* equals No_Element, then Constraint_Error is propagated; if *Position* does not designate an element in *Container*, then Program_Error is propagated. If an element equivalent to *New_Item* is already present in *Container* at a position other than *Position*, Program_Error is propagated. Otherwise, Replace_Element assigns *New_Item* to the element designated by *Position*. Any exception raised by the assignment is propagated.

35/2 **procedure** Query_Element
 (*Position* : **in** Cursor;
 Process : **not null access procedure** (Element : **in** Element_Type));

36/2 If *Position* equals No_Element, then Constraint_Error is propagated. Otherwise, Query_Element calls *Process.all* with the element designated by *Position* as the argument. Program_Error is propagated if *Process.all* tampers with the elements of *Container*. Any exception raised by *Process.all* is propagated.

37/2 **procedure** Move (*Target* : **in out** Set;
 Source : **in out** Set);

38/2 If *Target* denotes the same object as *Source*, then Move has no effect. Otherwise, Move first clears *Target*. Then, each element from *Source* is removed from *Source* and inserted into *Target*. The length of *Source* is 0 after a successful call to Move.

39/2 **procedure** Insert (*Container* : **in out** Set;
 New_Item : **in** Element_Type;
 Position : **out** Cursor;
 Inserted : **out** Boolean);

40/2 Insert checks if an element equivalent to *New_Item* is already present in *Container*. If a match is found, *Inserted* is set to False and *Position* designates the matching element. Otherwise, Insert adds *New_Item* to *Container*; *Inserted* is set to True and *Position* designates the newly-inserted element. Any exception raised during allocation is propagated and *Container* is not modified.

41/2 **procedure** Insert (*Container* : **in out** Set;
 New_Item : **in** Element_Type);

42/2 Insert inserts *New_Item* into *Container* as per the four-parameter Insert, with the difference that if an element equivalent to *New_Item* is already in the set, then Constraint_Error is propagated.

43/2 **procedure** Include (*Container* : **in out** Set;
 New_Item : **in** Element_Type);

44/2 Include inserts *New_Item* into *Container* as per the four-parameter Insert, with the difference that if an element equivalent to *New_Item* is already in the set, then it is replaced. Any exception raised during assignment is propagated.

45/2 **procedure** Replace (*Container* : **in out** Set;
 New_Item : **in** Element_Type);

46/2 Replace checks if an element equivalent to *New_Item* is already in the set. If a match is found, that element is replaced with *New_Item*; otherwise, Constraint_Error is propagated.

47/2 **procedure** Exclude (*Container* : **in out** Set;
 Item : **in** Element_Type);

48/2 Exclude checks if an element equivalent to *Item* is present in *Container*. If a match is found, Exclude removes the element from the set.

procedure Delete (Container : in out Set; Item : in Element_Type);	49/2
Delete checks if an element equivalent to Item is present in Container. If a match is found, Delete removes the element from the set; otherwise, Constraint_Error is propagated.	50/2
procedure Delete (Container : in out Set; Position : in out Cursor);	51/2
If Position equals No_Element, then Constraint_Error is propagated. If Position does not designate an element in Container, then Program_Error is propagated. Otherwise, Delete removes the element designated by Position from the set. Position is set to No_Element on return.	52/2
procedure Union (Target : in out Set; Source : in Set);	53/2
Union inserts into Target the elements of Source that are not equivalent to some element already in Target.	54/2
function Union (Left, Right : Set) return Set;	55/2
Returns a set comprising all of the elements of Left, and the elements of Right that are not equivalent to some element of Left.	56/2
procedure Intersection (Target : in out Set; Source : in Set);	57/2
Union deletes from Target the elements of Target that are not equivalent to some element of Source.	58/2
function Intersection (Left, Right : Set) return Set;	59/2
Returns a set comprising all the elements of Left that are equivalent to the some element of Right.	60/2
procedure Difference (Target : in out Set; Source : in Set);	61/2
If Target denotes the same object as Source, then Difference clears Target. Otherwise, it deletes from Target the elements that are equivalent to some element of Source.	62/2
function Difference (Left, Right : Set) return Set;	63/2
Returns a set comprising the elements of Left that are not equivalent to some element of Right.	64/2
procedure Symmetric_Difference (Target : in out Set; Source : in Set);	65/2
If Target denotes the same object as Source, then Symmetric_Difference clears Target. Otherwise, it deletes from Target the elements that are equivalent to some element of Source, and inserts into Target the elements of Source that are not equivalent to some element of Target.	66/2
function Symmetric_Difference (Left, Right : Set) return Set;	67/2
Returns a set comprising the elements of Left that are not equivalent to some element of Right, and the elements of Right that are not equivalent to some element of Left.	68/2
function Overlap (Left, Right : Set) return Boolean;	69/2
If an element of Left is equivalent to some element of Right, then Overlap returns True. Otherwise it returns False.	70/2

```

71/2      function Is_Subset (Subset : Set;
                           Of_Set : Set) return Boolean;
72/2      If an element of Subset is not equivalent to some element of Of_Set, then Is_Subset returns
              False. Otherwise it returns True.

73/2      function First (Container : Set) return Cursor;
74/2      If Length (Container) = 0, then First returns No_Element. Otherwise, First returns a cursor that
              designates the first element in Container.

75/2      function Next (Position : Cursor) return Cursor;
76/2      Returns a cursor that designates the successor of the element designated by Position. If Position
              designates the last element, then No_Element is returned. If Position equals No_Element, then
              No_Element is returned.

77/2      procedure Next (Position : in out Cursor);
78/2      Equivalent to Position := Next (Position).

79/2      Equivalent to Find (Container, Item) /= No_Element.

80/2      function Find (Container : Set;
                           Item       : Element_Type) return Cursor;
81/2      If Length (Container) equals 0, then Find returns No_Element. Otherwise, Find checks if an
              element equivalent to Item is present in Container. If a match is found, a cursor designating the
              matching element is returned; otherwise, No_Element is returned.

82/2      function Contains (Container : Set;
                           Item       : Element_Type) return Boolean;
83/2      function Has_Element (Position : Cursor) return Boolean;
84/2      Returns True if Position designates an element, and returns False otherwise.

85/2      procedure Iterate
              (Container : in Set;
               Process   : not null access procedure (Position : in Cursor));
86/2      Iterate calls Process.all with a cursor that designates each element in Container, starting with the
              first element and moving the cursor according to the successor relation. Program_Error is
              propagated if Process.all tampers with the cursors of Container. Any exception raised by
              Process.all is propagated.

87/2      Both Containers.Hashed_Set and Containers.Ordered_Set declare a nested generic package Generic_Keys,
              which provides operations that allow set manipulation in terms of a key (typically, a portion of an element)
              instead of a complete element. The formal function Key of Generic_Keys extracts a key value from an
              element. It is expected to return the same value each time it is called with a particular element. The
              behavior of Generic_Keys is unspecified if Key behaves in some other manner.

88/2      A key is expected to unambiguously determine a single equivalence class for elements. The behavior of
              Generic_Keys is unspecified if the formal parameters of this package behave in some other manner.

89/2      function Key (Position : Cursor) return Key_Type;
90/2      Equivalent to Key (Element (Position)).

```

The subprograms in package Generic_Keys named Contains, Find, Element, Delete, and Exclude, are equivalent to the corresponding subprograms in the parent package, with the difference that the Key parameter is used to locate an element in the set. 91/2

```
procedure Replace (Container : in out Set;
                  Key      : in      Key_Type;
                  New_Item : in      Element_Type); 92/2
```

Equivalent to Replace_Element (Container, Find (Container, Key), New_Item). 93/2

```
procedure Update_Element_Preserving_Key
  (Container : in out Set;
   Position  : in      Cursor;
   Process   : not null access procedure
              (Element : in out Element_Type)); 94/2
```

If Position equals No_Element, then Constraint_Error is propagated; if Position does not designate an element in Container, then Program_Error is propagated. Otherwise, Update_Element_Preserving_Key uses Key to save the key value K of the element designated by Position. Update_Element_Preserving_Key then calls Process.all with that element as the argument. Program_Error is propagated if Process.all tampers with the elements of Container. Any exception raised by Process.all is propagated. After Process.all returns, Update_Element_Preserving_Key checks if K determines the same equivalence class as that for the new element; if not, the element is removed from the set and Program_Error is propagated. 95/2

If Element_Type is unconstrained and definite, then the actual Element parameter of Process.all shall be unconstrained. 96/2

Erroneous Execution

A Cursor value is *invalid* if any of the following have occurred since it was created: 97/2

- The set that contains the element it designates has been finalized; 98/2
- The set that contains the element it designates has been used as the Source or Target of a call to Move; or 99/2
- The element it designates has been deleted from the set. 100/2

The result of "=" or Has_Element is unspecified if these functions are called with an invalid cursor parameter. Execution is erroneous if any other subprogram declared in Containers.Hashed_Sets or Containers.Ordered_Sets is called with an invalid cursor parameter. 101/2

Implementation Requirements

No storage associated with a Set object shall be lost upon assignment or scope exit. 102/2

The execution of an **assignment_statement** for a set shall have the effect of copying the elements from the source set object to the target set object. 103/2

Implementation Advice

Move should not copy elements, and should minimize copying of internal data structures. 104/2

If an exception is propagated from a set operation, no storage should be lost, nor any elements removed from a set unless specified by the operation. 105/2

A.18.8 The Package Containers.Hash_Sets

Static Semantics

1/2 The generic library package Containers.Hash_Sets has the following declaration:

```

2/2      generic
3/2          type Element_Type is private;
4/2          with function Hash (Element : Element_Type) return Hash_Type;
5/2          with function Equivalent_Elements (Left, Right : Element_Type)
6/2              return Boolean;
7/2          with function "=" (Left, Right : Element_Type) return Boolean is <>;
8/2      package Ada.Containers.Hash_Sets is
9/2          pragma Preelaborate(Hashed_Sets);

10/2         type Set is tagged private;
11/2         pragma Preelaborable_Initialization(Set);

12/2         type Cursor is private;
13/2         pragma Preelaborable_Initialization(Cursor);

14/2         Empty_Set : constant Set;
15/2         No_Element : constant Cursor;

16/2         function "=" (Left, Right : Set) return Boolean;
17/2         function Equivalent_Sets (Left, Right : Set) return Boolean;
18/2         function To_Set (New_Item : Element_Type) return Set;
19/2         function Capacity (Container : Set) return Count_Type;
20/2         procedure Reserve_Capacity (Container : in out Set;
21/2                                         Capacity : in      Count_Type);

22/2         function Length (Container : Set) return Count_Type;
23/2         function Is_Empty (Container : Set) return Boolean;
24/2         procedure Clear (Container : in out Set);
25/2         function Element (Position : Cursor) return Element_Type;
26/2         procedure Replace_Element (Container : in out Set;
27/2                                         Position : in      Cursor;
28/2                                         New_Item : in      Element_Type);

29/2         procedure Query_Element
30/2             (Position : in Cursor;
31/2              Process : not null access procedure (Element : in Element_Type));
32/2
33/2         procedure Move (Target : in out Set;
34/2                         Source : in out Set);

35/2         procedure Insert (Container : in out Set;
36/2                         New_Item : in      Element_Type;
37/2                         Position : out Cursor;
38/2                         Inserted : out Boolean);

39/2         procedure Insert (Container : in out Set;
40/2                         New_Item : in      Element_Type);

41/2         procedure Include (Container : in out Set;
42/2                         New_Item : in      Element_Type);

43/2         procedure Replace (Container : in out Set;
44/2                         New_Item : in      Element_Type);

45/2         procedure Exclude (Container : in out Set;
46/2                         Item : in      Element_Type);

47/2         procedure Delete (Container : in out Set;
48/2                         Item : in      Element_Type);

49/2         procedure Delete (Container : in out Set;
50/2                         Position : in out Cursor);

```

procedure Union (Target : in out Set;	Source : in Set);	26/2
function Union (Left, Right : Set) return Set;		27/2
function "or" (Left, Right : Set) return Set renames Union;		28/2
procedure Intersection (Target : in out Set;	Source : in Set);	29/2
function Intersection (Left, Right : Set) return Set;		30/2
function "and" (Left, Right : Set) return Set renames Intersection;		31/2
procedure Difference (Target : in out Set;	Source : in Set);	32/2
function Difference (Left, Right : Set) return Set;		33/2
function "-" (Left, Right : Set) return Set renames Difference;		34/2
procedure Symmetric_Difference (Target : in out Set;	Source : in Set);	35/2
function Symmetric_Difference (Left, Right : Set) return Set;		36/2
function "xor" (Left, Right : Set) return Set renames Symmetric_Difference;		37/2
function Overlap (Left, Right : Set) return Boolean;		38/2
function Is_Subset (Subset : Set;	Of_Set : Set) return Boolean;	39/2
function First (Container : Set) return Cursor;		40/2
function Next (Position : Cursor) return Cursor;		41/2
procedure Next (Position : in out Cursor);		42/2
function Find (Container : Set;	Item : Element_Type) return Cursor;	43/2
function Contains (Container : Set;	Item : Element_Type) return Boolean;	44/2
function Has_Element (Position : Cursor) return Boolean;		45/2
function Equivalent_Elements (Left, Right : Cursor)	return Boolean;	46/2
function Equivalent_Elements (Left : Cursor;	Right : Element_Type)	47/2
function Equivalent_Elements (Left : Element_Type;	Right : Cursor)	48/2
return Boolean;		
procedure Iterate	(Container : in Set;	49/2
Process : not null access procedure (Position : in Cursor));		
generic		50/2
type Key_Type (<>) is private;		
with function Key (Element : Element_Type) return Key_Type;		
with function Hash (Key : Key_Type) return Hash_Type;		
with function Equivalent_Keys (Left, Right : Key_Type)	return Boolean;	
package Generic_Keys is		
function Key (Position : Cursor) return Key_Type;		51/2
function Element (Container : Set;	Key : Key_Type)	52/2
return Element_Type;		
procedure Replace (Container : in out Set;		53/2
Key : in Key_Type;		
New_Item : in Element_Type);		

```

54/2      procedure Exclude (Container : in out Set;
                           Key      : in      Key_Type);
55/2      procedure Delete (Container : in out Set;
                           Key      : in      Key_Type);
56/2      function Find  (Container : Set;
                           Key      : Key_Type)
                           return Cursor;
57/2      function Contains (Container : Set;
                           Key      : Key_Type)
                           return Boolean;
58/2      procedure Update_Element_Preserving_Key
                           (Container : in out Set;
                            Position  : in      Cursor;
                            Process   : not null access procedure
                           (Element : in out Element_Type));
59/2      end Generic_Keys;
60/2      private
61/2          ... -- not specified by the language
62/2      end Ada.Containers.Hashed_Sets;

```

63/2 An object of type Set contains an expandable hash table, which is used to provide direct access to elements. The *capacity* of an object of type Set is the maximum number of elements that can be inserted into the hash table prior to it being automatically expanded.

- 64/2 Two elements *E1* and *E2* are defined to be *equivalent* if Equivalent_Elements (*E1*, *E2*) returns True.
- 65/2 The actual function for the generic formal function Hash is expected to return the same value each time it is called with a particular element value. For any two equivalent elements, the actual for Hash is expected to return the same value. If the actual for Hash behaves in some other manner, the behavior of this package is unspecified. Which subprograms of this package call Hash, and how many times they call it, is unspecified.
- 66/2 The actual function for the generic formal function Equivalent_Elements is expected to return the same value each time it is called with a particular pair of Element values. It should define an equivalence relationship, that is, be reflexive, symmetric, and transitive. If the actual for Equivalent_Elements behaves in some other manner, the behavior of this package is unspecified. Which subprograms of this package call Equivalent_Elements, and how many times they call it, is unspecified.
- 67/2 If the value of an element stored in a set is changed other than by an operation in this package such that at least one of Hash or Equivalent_Elements give different results, the behavior of this package is unspecified.
- 68/2 Which elements are the first element and the last element of a set, and which element is the successor of a given element, are unspecified, other than the general semantics described in A.18.7.

```

69/2      function Capacity (Container : Set) return Count_Type;
70/2          Returns the capacity of Container.
71/2      procedure Reserve_Capacity (Container : in out Set;
                                      Capacity  : in      Count_Type);
72/2          Reserve_Capacity allocates a new hash table such that the length of the resulting set can become
                      at least the value Capacity without requiring an additional call to Reserve_Capacity, and is large
                      enough to hold the current length of Container. Reserve_Capacity then rehashes the elements in
                      Container onto the new hash table. It replaces the old hash table with the new hash table, and

```

then deallocates the old hash table. Any exception raised during allocation is propagated and Container is not modified.

Reserve_Capacity tampers with the cursors of Container.

73/2

```
procedure Clear (Container : in out Set);
```

74/2

In addition to the semantics described in A.18.7, Clear does not affect the capacity of Container.

75/2

```
procedure Insert (Container : in out Set;
                  New_Item : in Element_Type;
                  Position : out Cursor;
                  Inserted : out Boolean);
```

76/2

In addition to the semantics described in A.18.7, if Length (Container) equals Capacity (Container), then Insert first calls Reserve_Capacity to increase the capacity of Container to some larger value.

77/2

```
function First (Container : Set) return Cursor;
```

78/2

If Length (Container) = 0, then First returns No_Element. Otherwise, First returns a cursor that designates the first hashed element in Container.

79/2

```
function Equivalent_Elements (Left, Right : Cursor)
                                return Boolean;
```

80/2

Equivalent to Equivalent_Elements (Element (Left), Element (Right)).

81/2

```
function Equivalent_Elements (Left : Cursor;
                                Right : Element_Type) return Boolean;
```

82/2

Equivalent to Equivalent_Elements (Element (Left), Right).

83/2

```
function Equivalent_Elements (Left : Element_Type;
                                Right : Cursor) return Boolean;
```

84/2

Equivalent to Equivalent_Elements (Left, Element (Right)).

85/2

For any element E , the actual function for the generic formal function Generic_Keys.Hash is expected to be such that Hash (E) = Generic_Keys.Hash (Key (E)). If the actuals for Key or Generic_Keys.Hash behave in some other manner, the behavior of Generic_Keys is unspecified. Which subprograms of Generic_Keys call Generic_Keys.Hash, and how many times they call it, is unspecified.

86/2

For any two elements $E1$ and $E2$, the boolean values Equivalent_Elements ($E1, E2$) and Equivalent_Keys (Key ($E1$), Key ($E2$)) are expected to be equal. If the actuals for Key or Equivalent_Keys behave in some other manner, the behavior of Generic_Keys is unspecified. Which subprograms of Generic_Keys call Equivalent_Keys, and how many times they call it, is unspecified.

87/2

Implementation Advice

If N is the length of a set, the average time complexity of the subprograms Insert, Include, Replace, Delete, Exclude and Find that take an element parameter should be $O(\log N)$. The average time complexity of the subprograms that take a cursor parameter should be $O(1)$. The average time complexity of Reserve_Capacity should be $O(N)$.

88/2

A.18.9 The Package Containers.Ordered_Sets

Static Semantics

1/2 The generic library package Containers.Ordered_Sets has the following declaration:

```

2/2   generic
      type Element_Type is private;
      with function "<" (Left, Right : Element_Type) return Boolean is <>;
      with function "=" (Left, Right : Element_Type) return Boolean is <>;
package Ada.Containers.Ordered_Sets is
  pragma Preelaborate(Ordered_Sets);
3/2   function Equivalent_Elements (Left, Right : Element_Type) return Boolean;
4/2   type Set is tagged private;
  pragma Preelaborable_Initialization(Set);
5/2   type Cursor is private;
  pragma Preelaborable_Initialization(Cursor);
6/2   Empty_Set : constant Set;
7/2   No_Element : constant Cursor;
8/2   function "=" (Left, Right : Set) return Boolean;
9/2   function Equivalent_Sets (Left, Right : Set) return Boolean;
10/2  function To_Set (New_Item : Element_Type) return Set;
11/2  function Length (Container : Set) return Count_Type;
12/2  function Is_Empty (Container : Set) return Boolean;
13/2  procedure Clear (Container : in out Set);
14/2  function Element (Position : Cursor) return Element_Type;
15/2  procedure Replace_Element (Container : in out Set;
                                Position : in      Cursor;
                                New_Item : in      Element_Type);
16/2  procedure Query_Element
      (Position : in Cursor;
       Process  : not null access procedure (Element : in Element_Type));
17/2  procedure Move (Target : in out Set;
                     Source : in out Set);
18/2  procedure Insert (Container : in out Set;
                        New_Item : in      Element_Type;
                        Position : out     Cursor;
                        Inserted : out     Boolean);
19/2  procedure Insert (Container : in out Set;
                        New_Item : in      Element_Type);
20/2  procedure Include (Container : in out Set;
                        New_Item : in      Element_Type);
21/2  procedure Replace (Container : in out Set;
                        New_Item : in      Element_Type);
22/2  procedure Exclude (Container : in out Set;
                        Item     : in      Element_Type);
23/2  procedure Delete (Container : in out Set;
                        Item     : in      Element_Type);
24/2  procedure Delete (Container : in out Set;
                        Position : in out Cursor);
25/2  procedure Delete_First (Container : in out Set);
26/2  procedure Delete_Last (Container : in out Set);

```

procedure Union (Target : in out Set;	27/2
Source : in Set);	
function Union (Left, Right : Set) return Set;	28/2
function "or" (Left, Right : Set) return Set renames Union;	29/2
procedure Intersection (Target : in out Set;	30/2
Source : in Set);	
function Intersection (Left, Right : Set) return Set;	31/2
function "and" (Left, Right : Set) return Set renames Intersection;	32/2
procedure Difference (Target : in out Set;	33/2
Source : in Set);	
function Difference (Left, Right : Set) return Set;	34/2
function "-" (Left, Right : Set) return Set renames Difference;	35/2
procedure Symmetric_Difference (Target : in out Set;	36/2
Source : in Set);	
function Symmetric_Difference (Left, Right : Set) return Set;	37/2
function "xor" (Left, Right : Set) return Set renames	38/2
Symmetric_Difference;	
function Overlap (Left, Right : Set) return Boolean;	39/2
function Is_Subset (Subset : Set;	40/2
Of_Set : Set) return Boolean;	
function First (Container : Set) return Cursor;	41/2
function First_Element (Container : Set) return Element_Type;	42/2
function Last (Container : Set) return Cursor;	43/2
function Last_Element (Container : Set) return Element_Type;	44/2
function Next (Position : Cursor) return Cursor;	45/2
procedure Next (Position : in out Cursor);	46/2
function Previous (Position : Cursor) return Cursor;	47/2
procedure Previous (Position : in out Cursor);	48/2
function Find (Container : Set;	49/2
Item : Element_Type)	
return Cursor;	
function Floor (Container : Set;	50/2
Item : Element_Type)	
return Cursor;	
function Ceiling (Container : Set;	51/2
Item : Element_Type)	
return Cursor;	
function Contains (Container : Set;	52/2
Item : Element_Type) return Boolean;	
function Has_Element (Position : Cursor) return Boolean;	53/2
function "<" (Left, Right : Cursor) return Boolean;	54/2
function ">" (Left, Right : Cursor) return Boolean;	55/2
function "<" (Left : Cursor; Right : Element_Type)	56/2
return Boolean;	
function ">" (Left : Cursor; Right : Element_Type)	57/2
return Boolean;	
function "<" (Left : Element_Type; Right : Cursor)	58/2
return Boolean;	
function ">" (Left : Element_Type; Right : Cursor)	59/2
return Boolean;	

```

60/2      procedure Iterate
              (Container : in Set;
               Process   : not null access procedure (Position : in Cursor));
61/2      procedure Reverse_Iterate
              (Container : in Set;
               Process   : not null access procedure (Position : in Cursor));
62/2      generic
              type Key_Type (<>) is private;
              with function Key (Element : Element_Type) return Key_Type;
              with function "<" (Left, Right : Key_Type)
                  return Boolean is <>;
              package Generic_Keys is
63/2          function Equivalent_Keys (Left, Right : Key_Type)
                          return Boolean;
64/2          function Key (Position : Cursor) return Key_Type;
65/2          function Element (Container : Set;
                                Key       : Key_Type)
                          return Element_Type;
66/2          procedure Replace (Container : in out Set;
                                Key      : in      Key_Type;
                                New_Item : in      Element_Type);
67/2          procedure Exclude (Container : in out Set;
                                Key      : in      Key_Type);
68/2          procedure Delete (Container : in out Set;
                                Key      : in      Key_Type);
69/2          function Find (Container : Set;
                            Key     : Key_Type)
                          return Cursor;
70/2          function Floor (Container : Set;
                            Key     : Key_Type)
                          return Cursor;
71/2          function Ceiling (Container : Set;
                                Key    : Key_Type)
                          return Cursor;
72/2          function Contains (Container : Set;
                                Key   : Key_Type) return Boolean;
73/2          procedure Update_Element_Preserving_Key
              (Container : in out Set;
               Position  : in      Cursor;
               Process   : not null access procedure
                           (Element : in out Element_Type));
74/2      end Generic_Keys;
75/2      private
76/2          ... -- not specified by the language
77/2      end Ada.Containers.Ordered_Sets;

```

78/2 Two elements $E1$ and $E2$ are *equivalent* if both $E1 < E2$ and $E2 < E1$ return False, using the generic formal " $<$ " operator for elements. Function Equivalent_Elements returns True if Left and Right are equivalent, and False otherwise.

79/2 The actual function for the generic formal function " $<$ " on Element_Type values is expected to return the same value each time it is called with a particular pair of key values. It should define a strict ordering relationship, that is, be irreflexive, asymmetric, and transitive. If the actual for " $<$ " behaves in some other manner, the behavior of this package is unspecified. Which subprograms of this package call " $<$ " and how many times they call it, is unspecified.

If the value of an element stored in a set is changed other than by an operation in this package such that at least one of "<" or "=" give different results, the behavior of this package is unspecified. 80/2

The first element of a nonempty set is the one which is less than all the other elements in the set. The last element of a nonempty set is the one which is greater than all the other elements in the set. The successor of an element is the smallest element that is larger than the given element. The predecessor of an element is the largest element that is smaller than the given element. All comparisons are done using the generic formal "<" operator for elements. 81/2

procedure Delete_First (Container : **in out** Set); 82/2

If Container is empty, Delete_First has no effect. Otherwise the element designated by First (Container) is removed from Container. Delete_First tampers with the cursors of Container. 83/2

procedure Delete_Last (Container : **in out** Set); 84/2

If Container is empty, Delete_Last has no effect. Otherwise the element designated by Last (Container) is removed from Container. Delete_Last tampers with the cursors of Container. 85/2

function First_Element (Container : Set) **return** Element_Type; 86/2

Equivalent to Element (First (Container)). 87/2

function Last (Container : Set) **return** Cursor; 88/2

Returns a cursor that designates the last element in Container. If Container is empty, returns No_Element. 89/2

function Last_Element (Container : Set) **return** Element_Type; 90/2

Equivalent to Element (Last (Container)). 91/2

function Previous (Position : Cursor) **return** Cursor; 92/2

If Position equals No_Element, then Previous returns No_Element. Otherwise Previous returns a cursor designating the element that precedes the one designated by Position. If Position designates the first element, then Previous returns No_Element. 93/2

procedure Previous (Position : **in out** Cursor); 94/2

Equivalent to Position := Previous (Position). 95/2

function Floor (Container : Set;
Item : Element_Type) **return** Cursor; 96/2

Floor searches for the last element which is not greater than Item. If such an element is found, a cursor that designates it is returned. Otherwise No_Element is returned. 97/2

function Ceiling (Container : Set;
Item : Element_Type) **return** Cursor; 98/2

Ceiling searches for the first element which is not less than Item. If such an element is found, a cursor that designates it is returned. Otherwise No_Element is returned. 99/2

function "<" (Left, Right : Cursor) **return** Boolean; 100/2

Equivalent to Element (Left) < Element (Right). 101/2

function ">" (Left, Right : Cursor) **return** Boolean; 102/2

Equivalent to Element (Right) < Element (Left). 103/2

```

104/2      function "<" (Left : Cursor; Right : Element_Type) return Boolean;
105/2          Equivalent to Element (Left) < Right.

106/2      function ">" (Left : Cursor; Right : Element_Type) return Boolean;
107/2          Equivalent to Right < Element (Left).

108/2      function "<" (Left : Element_Type; Right : Cursor) return Boolean;
109/2          Equivalent to Left < Element (Right).

110/2      function ">" (Left : Element_Type; Right : Cursor) return Boolean;
111/2          Equivalent to Element (Right) < Left.

112/2      procedure Reverse_Iterate
113/2          (Container : in Set;
113/2              Process   : not null access procedure (Position : in Cursor));
113/2          Iterates over the elements in Container as per Iterate, with the difference that the elements are
113/2          traversed in predecessor order, starting with the last element.

114/2  For any two elements  $E_1$  and  $E_2$ , the boolean values ( $E_1 < E_2$ ) and ( $\text{Key}(E_1) < \text{Key}(E_2)$ ) are expected to
114/2  be equal. If the actuals for Key or Generic_Keys." $<$ " behave in some other manner, the behavior of this
114/2  package is unspecified. Which subprograms of this package call Key and Generic_Keys." $<$ ", and how
114/2  many times the functions are called, is unspecified.

115/2  In addition to the semantics described in A.18.7, the subprograms in package Generic_Keys named Floor
115/2  and Ceiling, are equivalent to the corresponding subprograms in the parent package, with the difference
115/2  that the Key subprogram parameter is compared to elements in the container using the Key and " $<$ "
115/2  generic formal functions. The function named Equivalent_Keys in package Generic_Keys returns True if
115/2  both Left < Right and Right < Left return False using the generic formal " $<$ " operator, and returns True
115/2  otherwise.

```

Implementation Advice

116/2 If N is the length of a set, then the worst-case time complexity of the Insert, Include, Replace, Delete, Exclude and Find operations that take an element parameter should be $O(\log N)^{**}2$ or better. The worst-case time complexity of the subprograms that take a cursor parameter should be $O(1)$.

A.18.10 The Package Containers.Indefinite_Vectors

1/2 The language-defined generic package Containers.Indefinite_Vectors provides a private type Vector and a set of operations. It provides the same operations as the package Containers.Vectors (see A.18.2), with the difference that the generic formal Element_Type is indefinite.

Static Semantics

2/2 The declaration of the generic library package Containers.Indefinite_Vectors has the same contents as Containers.Vectors except:

- 3/2 • The generic formal Element_Type is indefinite.
- 4/2 • The procedures with the profiles:

5/2 **procedure** Insert (Container : **in out** Vector;
Before : **in** Extended_Index;
Count : **in** Count_Type := 1);

```
procedure Insert (Container : in out Vector;
                  Before   : in      Cursor;
                  Position  : out    Cursor;
                  Count    : in      Count_Type := 1);
```

are omitted.

- The actual Element parameter of access subprogram Process of Update_Element may be constrained even if Element_Type is unconstrained.

A.18.11 The Package Containers.Indefinite_Doubly_Linked_Lists

The language-defined generic package Containers.Indefinite_Doubly_Linked_Lists provides private types List and Cursor, and a set of operations for each type. It provides the same operations as the package Containers.Doubly_Linked_Lists (see A.18.3), with the difference that the generic formal Element_Type is indefinite.

Static Semantics

The declaration of the generic library package Containers.Indefinite_Doubly_Linked_Lists has the same contents as Containers.Doubly_Linked_Lists except:

- The generic formal Element_Type is indefinite.
- The procedure with the profile:

```
procedure Insert (Container : in out List;
                  Before   : in      Cursor;
                  Position  : out    Cursor;
                  Count    : in      Count_Type := 1);
```

is omitted.

- The actual Element parameter of access subprogram Process of Update_Element may be constrained even if Element_Type is unconstrained.

A.18.12 The Package Containers.Indefinite_Hashed_Maps

The language-defined generic package Containers.Indefinite_Hashed_Maps provides a map with the same operations as the package Containers.Hashed_Maps (see A.18.5), with the difference that the generic formal types Key_Type and Element_Type are indefinite.

Static Semantics

The declaration of the generic library package Containers.Indefinite_Hashed_Maps has the same contents as Containers.Hashed_Maps except:

- The generic formal Key_Type is indefinite.
- The generic formal Element_Type is indefinite.
- The procedure with the profile:

```
procedure Insert (Container : in out Map;
                  Key     : in      Key_Type;
                  Position : out    Cursor;
                  Inserted : out    Boolean);
```

is omitted.

- The actual Element parameter of access subprogram Process of Update_Element may be constrained even if Element_Type is unconstrained.

A.18.13 The Package Containers.Indefinite_Ordered_Maps

1/2 The language-defined generic package Containers.Indefinite_Ordered_Maps provides a map with the same operations as the package Containers.Ordered_Maps (see A.18.6), with the difference that the generic formal types Key_Type and Element_Type are indefinite.

Static Semantics

2/2 The declaration of the generic library package Containers.Indefinite_Ordered_Maps has the same contents as Containers.Ordered_Maps except:

- 3/2 • The generic formal Key_Type is indefinite.
- 4/2 • The generic formal Element_Type is indefinite.

5/2 • The procedure with the profile:

```
6/2   procedure Insert (Container : in out Map;
                      Key      : in   Key_Type;
                      Position  : out  Cursor;
                      Inserted  : out  Boolean);
```

7/2 is omitted.

- 8/2 • The actual Element parameter of access subprogram Process of Update_Element may be constrained even if Element_Type is unconstrained.

A.18.14 The Package Containers.Indefinite_Hashed_Sets

1/2 The language-defined generic package Containers.Indefinite_Hashed_Sets provides a set with the same operations as the package Containers.Hashed_Sets (see A.18.8), with the difference that the generic formal type Element_Type is indefinite.

Static Semantics

2/2 The declaration of the generic library package Containers.Indefinite_Hashed_Sets has the same contents as Containers.Hashed_Sets except:

- 3/2 • The generic formal Element_Type is indefinite.
- 4/2 • The actual Element parameter of access subprogram Process of Update_Element_-Preserving_Key may be constrained even if Element_Type is unconstrained.

A.18.15 The Package Containers.Indefinite_Ordered_Sets

1/2 The language-defined generic package Containers.Indefinite_Ordered_Sets provides a set with the same operations as the package Containers.Ordered_Sets (see A.18.9), with the difference that the generic formal type Element_Type is indefinite.

Static Semantics

2/2 The declaration of the generic library package Containers.Indefinite_Ordered_Sets has the same contents as Containers.Ordered_Sets except:

- 3/2 • The generic formal Element_Type is indefinite.
- 4/2 • The actual Element parameter of access subprogram Process of Update_Element_-Preserving_Key may be constrained even if Element_Type is unconstrained.

A.18.16 Array Sorting

The language-defined generic procedures `Containers.Generic_Array_Sort` and `Containers.Generic_Constrained_Array_Sort` provide sorting on arbitrary array types.

Static Semantics

The generic library procedure `Containers.Generic_Array_Sort` has the following declaration:

```
generic
  type Index_Type is (<>);
  type Element_Type is private;
  type Array_Type is array (Index_Type range <>) of Element_Type;
  with function "<" (Left, Right : Element_Type)
    return Boolean is <>;
procedure Ada.Containers.Generic_Array_Sort (Container : in out Array_Type);
pragma Pure(Ada.Containers.Generic_Array_Sort);
```

Reorders the elements of Container such that the elements are sorted smallest first as determined by the generic formal "<" operator provided. Any exception raised during evaluation of "<" is propagated.

The actual function for the generic formal function "<" of `Generic_Array_Sort` is expected to return the same value each time it is called with a particular pair of element values. It should define a strict ordering relationship, that is, be irreflexive, asymmetric, and transitive; it should not modify Container. If the actual for "<" behaves in some other manner, the behavior of the instance of `Generic_Array_Sort` is unspecified. How many times `Generic_Array_Sort` calls "<" is unspecified.

The generic library procedure `Containers.Generic_Constrained_Array_Sort` has the following declaration:

```
generic
  type Index_Type is (<>);
  type Element_Type is private;
  type Array_Type is array (Index_Type) of Element_Type;
  with function "<" (Left, Right : Element_Type)
    return Boolean is <>;
procedure Ada.Containers.Generic_Constrained_Array_Sort
  (Container : in out Array_Type);
pragma Pure(Ada.Containers.Generic_Constrained_Array_Sort);
```

Reorders the elements of Container such that the elements are sorted smallest first as determined by the generic formal "<" operator provided. Any exception raised during evaluation of "<" is propagated.

The actual function for the generic formal function "<" of `Generic_Constrained_Array_Sort` is expected to return the same value each time it is called with a particular pair of element values. It should define a strict ordering relationship, that is, be irreflexive, asymmetric, and transitive; it should not modify Container. If the actual for "<" behaves in some other manner, the behavior of the instance of `Generic_Constrained_Array_Sort` is unspecified. How many times `Generic_Constrained_Array_Sort` calls "<" is unspecified.

Implementation Advice

The worst-case time complexity of a call on an instance of `Containers.Generic_Array_Sort` or `Containers.Generic_Constrained_Array_Sort` should be $O(N^{**}2)$ or better, and the average time complexity should be better than $O(N^{**}2)$, where N is the length of the Container parameter.

- 11/2 Containers.Generic_Array_Sort and Containers.Generic_Constrained_Array_Sort should minimize copying of elements.

Annex B (normative) Interface to Other Languages

This Annex describes features for writing mixed-language programs. General interface support is presented first; then specific support for C, COBOL, and Fortran is defined, in terms of language interface packages for each of these languages.

B.1 Interfacing Pragmas

A **pragma Import** is used to import an entity defined in a foreign language into an Ada program, thus allowing a foreign-language subprogram to be called from Ada, or a foreign-language variable to be accessed from Ada. In contrast, a **pragma Export** is used to export an Ada entity to a foreign language, thus allowing an Ada subprogram to be called from a foreign language, or an Ada object to be accessed from a foreign language. The **pragmas Import** and **Export** are intended primarily for objects and subprograms, although implementations are allowed to support other entities.

A **pragma Convention** is used to specify that an Ada entity should use the conventions of another language. It is intended primarily for types and “callback” subprograms. For example, “**pragma Convention(Fortran, Matrix);**” implies that Matrix should be represented according to the conventions of the supported Fortran implementation, namely column-major order.

A **pragma Linker_Options** is used to specify the system linker parameters needed when a given compilation unit is included in a partition.

Syntax

An *interfacing pragma* is a representation **pragma** that is one of the **pragmas Import**, **Export**, or **Convention**. Their forms, together with that of the related **pragma Linker_Options**, are as follows:

```

pragma Import
  [Convention =>] convention_identifier, [Entity =>] local_name
  [, [External_Name =>] string_expression] [, [Link_Name =>] string_expression]];

pragma Export
  [Convention =>] convention_identifier, [Entity =>] local_name
  [, [External_Name =>] string_expression] [, [Link_Name =>] string_expression]];

pragma Convention([Convention =>] convention_identifier, [Entity =>] local_name);

pragma Linker_Options(string_expression);

```

A **pragma Linker_Options** is allowed only at the place of a **declarative_item**.

For **pragmas Import** and **Export**, the argument for **Link_Name** shall not be given without the **pragma_argument_identifier** unless the argument for **External_Name** is given.

Name Resolution Rules

The expected type for a *string_expression* in an interfacing pragma or in **pragma Linker_Options** is **String**.

Legality Rules

The *convention_identifier* of an interfacing pragma shall be the name of a *convention*. The convention names are implementation defined, except for certain language-defined ones, such as Ada and Intrinsic, as

explained in 6.3.1, “Conformance Rules”. Additional convention names generally represent the calling conventions of foreign languages, language implementations, or specific run-time models. The convention of a callable entity is its *calling convention*.

If L is a *convention_identifier* for a language, then a type T is said to be *compatible with convention L*, (alternatively, is said to be an *L-compatible type*) if any of the following conditions are met:

- T is declared in a language interface package corresponding to L and is defined to be L -compatible (see B.3, B.3.1, B.3.2, B.4, B.5),
- Convention L has been specified for T in a **pragma Convention**, and T is *eligible for convention L*; that is:
 - T is an array type with either an unconstrained or statically-constrained first subtype, and its component type is L -compatible,
 - T is a record type that has no discriminants and that only has components with statically-constrained subtypes, and each component type is L -compatible,
 - T is an access-to-object type, and its designated type is L -compatible,
 - T is an access-to-subprogram type, and its designated profile's parameter and result types are all L -compatible.
- T is derived from an L -compatible type,
- The implementation permits T as an L -compatible type.

If **pragma Convention** applies to a type, then the type shall either be compatible with or eligible for the convention specified in the **pragma**.

A **pragma Import** shall be the completion of a declaration. Notwithstanding any rule to the contrary, a **pragma Import** may serve as the completion of any kind of (explicit) declaration if supported by an implementation for that kind of declaration. If a completion is a **pragma Import**, then it shall appear in the same **declarative_part**, **package_specification**, **task_definition** or **protected_definition** as the declaration. For a library unit, it shall appear in the same compilation, before any subsequent **compilation_units** other than **pragmas**. If the **local_name** denotes more than one entity, then the **pragma Import** is the completion of all of them.

An entity specified as the **Entity** argument to a **pragma Import** (or **pragma Export**) is said to be *imported* (respectively, *exported*).

The declaration of an imported object shall not include an explicit initialization expression. Default initializations are not performed.

The type of an imported or exported object shall be compatible with the convention specified in the corresponding **pragma**.

For an imported or exported subprogram, the result and parameter types shall each be compatible with the convention specified in the corresponding **pragma**.

The external name and link name **string_expressions** of a **pragma Import** or **Export**, and the **string_expression** of a **pragma Linker_Options**, shall be static.

Static Semantics

Import, **Export**, and **Convention pragmas** are representation pragmas that specify the *convention* aspect of representation. In addition, **Import** and **Export pragmas** specify the *imported* and *exported* aspects of representation, respectively.

An interfacing pragma is a program unit pragma when applied to a program unit (see 10.1.5). 29
 An interfacing pragma defines the convention of the entity denoted by the `local_name`. The convention 30
 represents the calling convention or representation convention of the entity. For an access-to-subprogram
 type, it represents the calling convention of designated subprograms. In addition:

- A `pragma Import` specifies that the entity is defined externally (that is, outside the Ada 31
 program).
- A `pragma Export` specifies that the entity is used externally. 32
- A `pragma Import` or `Export` optionally specifies an entity's external name, link name, or both. 33

An *external name* is a string value for the name used by a foreign language program either for an entity 34
 that an Ada program imports, or for referring to an entity that an Ada program exports.

A *link name* is a string value for the name of an exported or imported entity, based on the conventions of 35
 the foreign language's compiler in interfacing with the system's linker tool.

The meaning of link names is implementation defined. If neither a link name nor the `Address` attribute of 36
 an imported or exported entity is specified, then a link name is chosen in an implementation-defined
 manner, based on the external name if one is specified.

`Pragma Linker_Options` has the effect of passing its string argument as a parameter to the system linker (if 37
 one exists), if the immediately enclosing compilation unit is included in the partition being linked. The
 interpretation of the string argument, and the way in which the string arguments from multiple
`Linker_Options` pragmas are combined, is implementation defined.

Dynamic Semantics

Notwithstanding what this International Standard says elsewhere, the elaboration of a declaration denoted 38
 by the `local_name` of a `pragma Import` does not create the entity. Such an elaboration has no other effect
 than to allow the defining name to denote the external entity.

Erroneous Execution

It is the programmer's responsibility to ensure that the use of interfacing pragmas does not violate Ada 38.1/2
 semantics; otherwise, program execution is erroneous.

Implementation Advice

If an implementation supports `pragma Export` to a given language, then it should also allow the main 39
 subprogram to be written in that language. It should support some mechanism for invoking the elaboration
 of the Ada library units included in the system, and for invoking the finalization of the environment task.
 On typical systems, the recommended mechanism is to provide two subprograms whose link names are
 "adainit" and "adafinal". `Adainit` should contain the elaboration code for library units. `Adafinal` should
 contain the finalization code. These subprograms should have no effect the second and subsequent time
 they are called.

Automatic elaboration of preelaborated packages should be provided when `pragma Export` is supported. 40

For each supported convention *L* other than *Intrinsic*, an implementation should support `Import` and `Export` 41
 pragmas for objects of *L*-compatible types and for subprograms, and `pragma Convention` for *L*-eligible
 types and for subprograms, presuming the other language has corresponding features. `Pragma Convention`
 need not be supported for scalar types.

NOTES

- 42 1 Implementations may place restrictions on interfacing pragmas; for example, requiring each exported entity to be declared at the library level.
- 43 2 A pragma Import specifies the conventions for accessing external entities. It is possible that the actual entity is written in assembly language, but reflects the conventions of a particular language. For example, **pragma** Import(Ada, ...) can be used to interface to an assembly language routine that obeys the Ada compiler's calling conventions.
- 44 3 To obtain "call-back" to an Ada subprogram from a foreign language environment, **pragma** Convention should be specified both for the access-to-subprogram type and the specific subprogram(s) to which 'Access is applied.
- 45 4 It is illegal to specify more than one of Import, Export, or Convention for a given entity.
- 46 5 The local_name in an interfacing pragma can denote more than one entity in the case of overloading. Such a pragma applies to all of the denoted entities.
- 47 6 See also 13.8, "Machine Code Insertions".
- 48 7 If both External_Name and Link_Name are specified for an Import or Export pragma, then the External_Name is ignored.
- 49/2 *This paragraph was deleted.*

Examples

- 50 *Example of interfacing pragmas:*

```
51    package Fortran_Library is
      function Sqrt (X : Float) return Float;
      function Exp  (X : Float) return Float;
    private
      pragma Import(Fortran, Sqrt);
      pragma Import(Fortran, Exp);
    end Fortran_Library;
```

B.2 The Package Interfaces

- 1 Package Interfaces is the parent of several library packages that declare types and other entities useful for interfacing to foreign languages. It also contains some implementation-defined types that are useful across more than one language (in particular for interfacing to assembly language).

Static Semantics

- 2 The library package Interfaces has the following skeletal declaration:

```
3    package Interfaces is
      pragma Pure(Interfaces);
      type Integer_n is range -2** (n-1) .. 2** (n-1) - 1; --2's complement
      type Unsigned_n is mod 2**n;
      function Shift_Left  (Value : Unsigned_n; Amount : Natural)
        return Unsigned_n;
      function Shift_Right (Value : Unsigned_n; Amount : Natural)
        return Unsigned_n;
      function Shift_Right_Arithmetic (Value : Unsigned_n; Amount : Natural)
        return Unsigned_n;
      function Rotate_Left  (Value : Unsigned_n; Amount : Natural)
        return Unsigned_n;
      function Rotate_Right (Value : Unsigned_n; Amount : Natural)
        return Unsigned_n;
      ...
    end Interfaces;
```

Implementation Requirements

- An implementation shall provide the following declarations in the visible part of package Interfaces: 7
- Signed and modular integer types of n bits, if supported by the target architecture, for each n that is at least the size of a storage element and that is a factor of the word size. The names of these types are of the form Integer_ n for the signed types, and Unsigned_ n for the modular types; 8
 - For each such modular type in Interfaces, shifting and rotating subprograms as specified in the declaration of Interfaces above. These subprograms are Intrinsic. They operate on a bit-by-bit basis, using the binary representation of the value of the operands to yield a binary representation for the result. The Amount parameter gives the number of bits by which to shift or rotate. For shifting, zero bits are shifted in, except in the case of Shift_Right_Arithmetic, where one bits are shifted in if Value is at least half the modulus; 9
 - Floating point types corresponding to each floating point format fully supported by the hardware. 10

Support for interfacing to any foreign language is optional. However, an implementation shall not provide any attribute, library unit, or pragma having the same name as an attribute, library unit, or pragma (respectively) specified in the following clauses of this Annex unless the provided construct is either as specified in those clauses or is more limited in capability than that required by those clauses. A program that attempts to use an unsupported capability of this Annex shall either be identified by the implementation before run time or shall raise an exception at run time. 10.1/2

Implementation Permissions

An implementation may provide implementation-defined library units that are children of Interfaces, and may add declarations to the visible part of Interfaces in addition to the ones defined above. 11

A child package of package Interfaces with the name of a convention may be provided independently of whether the convention is supported by the pragma Convention and vice versa. Such a child package should contain any declarations that would be useful for interfacing to the language (implementation) represented by the convention. Any declarations useful for interfacing to any language on the given hardware architecture should be provided directly in Interfaces. 11.1/2

Implementation Advice

This paragraph was deleted. 12/2

An implementation supporting an interface to C, COBOL, or Fortran should provide the corresponding package or packages described in the following clauses. 13

B.3 Interfacing with C and C++

The facilities relevant to interfacing with the C language and the corresponding subset of the C++ language are the package Interfaces.C and its children; support for the Import, Export, and Convention pragmas with *convention_identifier* C; and support for the Convention pragma with *convention_identifier* C_Pass_By_Copy. 1/2

The package Interfaces.C contains the basic types, constants and subprograms that allow an Ada program to pass scalars and strings to C and C++ functions. When this clause mentions a C entity, the reference also applies to the corresponding entity in C++. 2/2

Static Semantics

3 The library package Interfaces.C has the following declaration:

```

4   package Interfaces.C is
5     pragma Pure(C);
6     -- Declarations based on C's <limits.h>
7     CHAR_BIT : constant := implementation-defined; -- typically 8
8     SCHAR_MIN : constant := implementation-defined; -- typically -128
9     SCHAR_MAX : constant := implementation-defined; -- typically 127
10    UCHAR_MAX : constant := implementation-defined; -- typically 255
11
12    -- Signed and Unsigned Integers
13    type int is range implementation-defined;
14    type short is range implementation-defined;
15    type long is range implementation-defined;
16
17    type signed_char is range SCHAR_MIN .. SCHAR_MAX;
18    for signed_char'Size use CHAR_BIT;
19
20    type unsigned is mod implementation-defined;
21    type unsigned_short is mod implementation-defined;
22    type unsigned_long is mod implementation-defined;
23
24    type unsigned_char is mod (UCHAR_MAX+1);
25    for unsigned_char'Size use CHAR_BIT;
26
27    subtype plain_char is implementation-defined;
28    type ptrdiff_t is range implementation-defined;
29    type size_t is mod implementation-defined;
30
31    -- Floating Point
32    type C_float      is digits implementation-defined;
33    type double       is digits implementation-defined;
34    type long_double  is digits implementation-defined;
35
36    -- Characters and Strings
37    type char is <implementation-defined character type>;
38    nul : constant char := implementation-defined;
39
40    function To_C (Item : in Character) return char;
41    function To_Ada (Item : in char) return Character;
42
43    type char_array is array (size_t range <>) of aliased char;
44    pragma Pack(char_array);
45    for char_array'Component_Size use CHAR_BIT;
46
47    function Is_Nul_Terminated (Item : in char_array) return Boolean;
48
49    function To_C (Item      : in String;
50                  Append_Nul : in Boolean := True)
51      return char_array;
52
53    function To_Ada (Item      : in char_array;
54                      Trim_Nul : in Boolean := True)
55      return String;
56
56    procedure To_C (Item      : in String;
57                    Target    : out char_array;
58                    Count    : out size_t;
59                    Append_Nul : in Boolean := True);
59
60    procedure To_Ada (Item      : in char_array;
61                      Target    : out String;
62                      Count    : out Natural;
63                      Trim_Nul : in Boolean := True);
64
65
66    -- Wide Character and Wide String
67    type wchar_t is <implementation-defined character type>;

```

```

wide_nul : constant wchar_t := implementation-defined; 31/1
function To_C (Item : in Wide_Character) return wchar_t; 32
function To_Ada (Item : in wchar_t) return Wide_Character;
type wchar_array is array (size_t range <>) of aliased wchar_t; 33
pragma Pack(wchar_array);
function Is_Nul_Terminated (Item : in wchar_array) return Boolean; 35
function To_C (Item : in Wide_String;
               Append_Nul : in Boolean := True)
    return wchar_array; 36
function To_Ada (Item : in wchar_array;
                 Trim_Nul : in Boolean := True) 37
    return Wide_String;
procedure To_C (Item : in Wide_String;
               Target : out wchar_array;
               Count : out size_t;
               Append_Nul : in Boolean := True); 38
procedure To_Ada (Item : in wchar_array;
                  Target : out Wide_String;
                  Count : out Natural;
                  Trim_Nul : in Boolean := True); 39
-- ISO/IEC 10646:2003 compatible types defined by ISO/IEC TR 19769:2004. 39.1/2
type char16_t is <implementation-defined character type>; 39.2/2
char16_nul : constant char16_t := implementation-defined; 39.3/2
function To_C (Item : in Wide_Character) return char16_t; 39.4/2
function To_Ada (Item : in char16_t) return Wide_Character;
type char16_array is array (size_t range <>) of aliased char16_t; 39.5/2
pragma Pack(char16_array); 39.6/2
function Is_Nul_Terminated (Item : in char16_array) return Boolean; 39.7/2
function To_C (Item : in Wide_String;
               Append_Nul : in Boolean := True)
    return char16_array; 39.8/2
function To_Ada (Item : in char16_array;
                 Trim_Nul : in Boolean := True)
    return Wide_String;
procedure To_C (Item : in Wide_String;
               Target : out char16_array;
               Count : out size_t;
               Append_Nul : in Boolean := True); 39.9/2
procedure To_Ada (Item : in char16_array;
                  Target : out Wide_String;
                  Count : out Natural;
                  Trim_Nul : in Boolean := True); 39.10/2
type char32_t is <implementation-defined character type>; 39.11/2
char32_nul : constant char32_t := implementation-defined; 39.12/2
function To_C (Item : in Wide_Wide_Character) return char32_t; 39.13/2
function To_Ada (Item : in char32_t) return Wide_Wide_Character;
type char32_array is array (size_t range <>) of aliased char32_t; 39.14/2
pragma Pack(char32_array);
function Is_Nul_Terminated (Item : in char32_array) return Boolean; 39.16/2
function To_C (Item : in Wide_Wide_String;
               Append_Nul : in Boolean := True)
    return char32_array; 39.17/2
function To_Ada (Item : in char32_array;
                 Trim_Nul : in Boolean := True)
    return Wide_Wide_String;

```

```

39.18/2      procedure To_C (Item      : in Wide_Wide_String;
                           Target    : out char32_array;
                           Count     : out size_t;
                           Append_Nul : in Boolean := True);
39.19/2      procedure To_Ada (Item      : in char32_array;
                           Target    : out Wide_Wide_String;
                           Count     : out Natural;
                           Trim_Nul  : in Boolean := True);
40          Terminator_Error : exception;
41      end Interfaces.C;

```

42 Each of the types declared in Interfaces.C is C-compatible.

43/2 The types int, short, long, unsigned, ptrdiff_t, size_t, double, char, wchar_t, char16_t, and char32_t correspond respectively to the C types having the same names. The types signed_char, unsigned_short, unsigned_long, unsigned_char, C_float, and long_double correspond respectively to the C types signed char, unsigned short, unsigned long, unsigned char, float, and long double.

44 The type of the subtype plain_char is either signed_char or unsigned_char, depending on the C implementation.

```

45      function To_C   (Item : in Character) return char;
        function To_Ada (Item : in char)      return Character;

```

46 The functions To_C and To_Ada map between the Ada type Character and the C type char.

```

47      function Is_Nul_Terminated (Item : in char_array) return Boolean;

```

48 The result of Is_Nul_Terminated is True if Item contains nul, and is False otherwise.

```

49      function To_C   (Item : in String;      Append_Nul : in Boolean := True)
                    return char_array;
      function To_Ada (Item : in char_array; Trim_Nul   : in Boolean := True)
                    return String;

```

50/2 The result of To_C is a char_array value of length Item'Length (if Append_Nul is False) or Item'Length+1 (if Append_Nul is True). The lower bound is 0. For each component Item(I), the corresponding component in the result is To_C applied to Item(I). The value nul is appended if Append_Nul is True. If Append_Nul is False and Item'Length is 0, then To_C propagates Constraint_Error.

51 The result of To_Ada is a String whose length is Item'Length (if Trim_Nul is False) or the length of the slice of Item preceding the first nul (if Trim_Nul is True). The lower bound of the result is 1. If Trim_Nul is False, then for each component Item(I) the corresponding component in the result is To_Ada applied to Item(I). If Trim_Nul is True, then for each component Item(I) before the first nul the corresponding component in the result is To_Ada applied to Item(I). The function propagates Terminator_Error if Trim_Nul is True and Item does not contain nul.

```
procedure To_C (Item      : in String;
                Target    : out char_array;
                Count     : out size_t;
                Append_Nul : in Boolean := True);
```

52

```
procedure To_Ada (Item      : in char_array;
                  Target    : out String;
                  Count     : out Natural;
                  Trim_Nul : in Boolean := True);
```

53

For procedure To_C, each element of Item is converted (via the To_C function) to a char, which is assigned to the corresponding element of Target. If Append_Nul is True, nul is then assigned to the next element of Target. In either case, Count is set to the number of Target elements assigned. If Target is not long enough, Constraint_Error is propagated.

For procedure To_Ada, each element of Item (if Trim_Nul is False) or each element of Item preceding the first nul (if Trim_Nul is True) is converted (via the To_Ada function) to a Character, which is assigned to the corresponding element of Target. Count is set to the number of Target elements assigned. If Target is not long enough, Constraint_Error is propagated. If Trim_Nul is True and Item does not contain nul, then Terminator_Error is propagated.

54

```
function Is_Nul_Terminated (Item : in wchar_array) return Boolean;
```

55

The result of Is_Nul_Terminated is True if Item contains wide_nul, and is False otherwise.

56

```
function To_C (Item : in Wide_Character) return wchar_t;
function To_Ada (Item : in wchar_t) return Wide_Character;
```

57

To_C and To_Ada provide the mappings between the Ada and C wide character types.

58

```
function To_C (Item      : in Wide_String;
                Append_Nul : in Boolean := True)
return wchar_array;
```

59

```
function To_Ada (Item      : in wchar_array;
                  Trim_Nul : in Boolean := True)
return Wide_String;
```

```
procedure To_C (Item      : in Wide_String;
                Target    : out wchar_array;
                Count     : out size_t;
                Append_Nul : in Boolean := True);
```

```
procedure To_Ada (Item      : in wchar_array;
                  Target    : out Wide_String;
                  Count     : out Natural;
                  Trim_Nul : in Boolean := True);
```

The To_C and To_Ada subprograms that convert between Wide_String and wchar_array have analogous effects to the To_C and To_Ada subprograms that convert between String and char_array, except that wide_nul is used instead of nul.

60

```
function Is_Nul_Terminated (Item : in char16_array) return Boolean;
```

60.1/2

The result of Is_Nul_Terminated is True if Item contains char16_nul, and is False otherwise.

60.2/2

```
function To_C (Item : in Wide_Character) return char16_t;
function To_Ada (Item : in char16_t) return Wide_Character;
```

60.3/2

To_C and To_Ada provide mappings between the Ada and C 16-bit character types.

60.4/2

```

60.5/2      function To_C (Item      : in Wide_String;
                      Append_Nul : in Boolean := True)
            return char16_array;

      function To_Ada (Item      : in char16_array;
                       Trim_Nul : in Boolean := True)
            return Wide_String;

      procedure To_C (Item      : in Wide_String;
                      Target    : out char16_array;
                      Count     : out size_t;
                      Append_Nul : in Boolean := True);

      procedure To_Ada (Item      : in char16_array;
                        Target    : out Wide_String;
                        Count     : out Natural;
                        Trim_Nul : in Boolean := True);

```

60.6/2 The To_C and To_Ada subprograms that convert between Wide_String and char16_array have analogous effects to the To_C and To_Ada subprograms that convert between String and char_array, except that char16_nul is used instead of nul.

60.7/2 **function** Is_Nul_Terminated (Item : in char32_array) **return** Boolean;

60.8/2 The result of Is_Nul_Terminated is True if Item contains char16_nul, and is False otherwise.

60.9/2 **function** To_C (Item : in Wide_Wide_Character) **return** char32_t;
function To_Ada (Item : in char32_t) **return** Wide_Wide_Character;

60.10/2 To_C and To_Ada provide mappings between the Ada and C 32-bit character types.

60.11/2 **function** To_C (Item : in Wide_Wide_String;
 Append_Nul : in Boolean := True)
 return char32_array;

function To_Ada (Item : in char32_array;
 Trim_Nul : in Boolean := True)
 return Wide_Wide_String;

procedure To_C (Item : in Wide_Wide_String;
 Target : out char32_array;
 Count : out size_t;
 Append_Nul : in Boolean := True);

procedure To_Ada (Item : in char32_array;
 Target : out Wide_Wide_String;
 Count : out Natural;
 Trim_Nul : in Boolean := True);

60.12/2 The To_C and To_Ada subprograms that convert between Wide_Wide_String and char32_array have analogous effects to the To_C and To_Ada subprograms that convert between String and char_array, except that char32_nul is used instead of nul.

60.13/1 A Convention pragma with *convention_identifier* C_Pass_By_Copy shall only be applied to a type.

60.14/2 The eligibility rules in B.1 do not apply to convention C_Pass_By_Copy. Instead, a type T is eligible for convention C_Pass_By_Copy if T is an unchecked union type or if T is a record type that has no discriminants and that only has components with statically constrained subtypes, and each component is C-compatible.

60.15/1 If a type is C_Pass_By_Copy-compatible then it is also C-compatible.

Implementation Requirements

An implementation shall support pragma Convention with a C *convention_identifier* for a C-eligible type (see B.1). An implementation shall support pragma Convention with a C_Pass_By_Copy *convention_identifier* for a C_Pass_By_Copy-eligible type. 61/1

Implementation Permissions

An implementation may provide additional declarations in the C interface packages. 62

Implementation Advice

The constants nul, wide_nul, char16_nul, and char32_nul should have a representation of zero. 62.1/2

An implementation should support the following interface correspondences between Ada and C. 63

- An Ada procedure corresponds to a void-returning C function. 64
- An Ada function corresponds to a non-void C function. 65
- An Ada **in** scalar parameter is passed as a scalar argument to a C function. 66
- An Ada **in** parameter of an access-to-object type with designated type T is passed as a t* argument to a C function, where t is the C type corresponding to the Ada type T. 67
- An Ada **access** T parameter, or an Ada **out** or **in out** parameter of an elementary type T, is passed as a t* argument to a C function, where t is the C type corresponding to the Ada type T. In the case of an elementary **out** or **in out** parameter, a pointer to a temporary copy is used to preserve by-copy semantics. 68
- An Ada parameter of a (record) type T of convention C_Pass_By_Copy, of mode **in**, is passed as a t argument to a C function, where t is the C struct corresponding to the Ada type T. 68.1/2
- An Ada parameter of a record type T, of any mode, other than an **in** parameter of a type of convention C_Pass_By_Copy, is passed as a t* argument to a C function, where t is the C struct corresponding to the Ada type T. 69/2
- An Ada parameter of an array type with component type T, of any mode, is passed as a t* argument to a C function, where t is the C type corresponding to the Ada type T. 70
- An Ada parameter of an access-to-subprogram type is passed as a pointer to a C function whose prototype corresponds to the designated subprogram's specification. 71

An Ada parameter of a private type is passed as specified for the full view of the type. 71.1/2

NOTES

8 Values of type char_array are not implicitly terminated with nul. If a char_array is to be passed as a parameter to an imported C function requiring nul termination, it is the programmer's responsibility to obtain this effect. 72

9 To obtain the effect of C's sizeof(item_type), where Item_Type is the corresponding Ada type, evaluate the expression: size_t(Item_Type'Size/CHAR_BIT). 73

This paragraph was deleted. 74/2

10 A C function that takes a variable number of arguments can correspond to several Ada subprograms, taking various specific numbers and types of parameters. 75

Examples

```

76  Example of using the Interfaces.C package:
77      --Calling the C Library Function strcpy
    with Interfaces.C;
    procedure Test is
        package C renames Interfaces.C;
        use type C.char_array;
        -- Call <string.h>strcpy:
        -- C definition of strcpy: char *strcpy(char *s1, const char *s2);
        -- This function copies the string pointed to by s2 (including the terminating null character)
        -- into the array pointed to by s1. If copying takes place between objects that overlap,
        -- the behavior is undefined. The strcpy function returns the value of s1.
78      -- Note: since the C function's return value is of no interest, the Ada interface is a procedure
        procedure Strcpy (Target : out C.char_array;
                           Source : in C.char_array);

79      pragma Import(C, Strcpy, "strcpy");
80      Chars1 : C.char_array(1..20);
81      Chars2 : C.char_array(1..20);
82
83      begin
84          Chars2(1..6) := "qwert" & C.nul;
85          Strcpy(Chars1, Chars2);
86          -- Now Chars1(1..6) = "qwert" & C.Nul
87      end Test;

```

B.3.1 The Package Interfaces.C.Strings

1 The package Interfaces.C.Strings declares types and subprograms allowing an Ada program to allocate, reference, update, and free C-style strings. In particular, the private type chars_ptr corresponds to a common use of “char **” in C programs, and an object of this type can be passed to a subprogram to which pragma Import(C,...) has been applied, and for which “char **” is the type of the argument of the C function.

Static Semantics

2 The library package Interfaces.C.Strings has the following declaration:

```

3  package Interfaces.C.Strings is
4      pragma Preelaborate(Strings);
5
6      type char_array_access is access all char_array;
7
8      type chars_ptr is private;
9      pragma Preelaborable_Initialization(chars_ptr);
10
11      type chars_ptr_array is array (size_t range <>) of aliased chars_ptr;
12      Null_Ptr : constant chars_ptr;
13
14      function To_Chars_Ptr (Item       : in char_array_access;
15                             Nul_Check : in Boolean := False)
16          return chars_ptr;
17
18      function New_Char_Array (Chars     : in char_array) return chars_ptr;
19
20      function New_String (Str : in String) return chars_ptr;
21
22      procedure Free (Item : in out chars_ptr);
23
24      Dereference_Error : exception;
25
26      function Value (Item : in chars_ptr) return char_array;
27
28      function Value (Item : in chars_ptr; Length : in size_t)
29          return char_array;
30
31      function Value (Item : in chars_ptr) return String;

```

```

function Value (Item : in chars_ptr; Length : in size_t) 16
    return String;
function Strlen (Item : in chars_ptr) return size_t; 17
procedure Update (Item : in chars_ptr; 18
                  Offset : in size_t;
                  Chars : in char_array;
                  Check : in Boolean := True);
procedure Update (Item : in chars_ptr; 19
                  Offset : in size_t;
                  Str : in String;
                  Check : in Boolean := True);
Update_Error : exception; 20
private 21
... -- not specified by the language
end Interfaces.C.Strings;

```

The type chars_ptr is C-compatible and corresponds to the use of C's "char *" for a pointer to the first char in a char array terminated by nul. When an object of type chars_ptr is declared, its value is by default set to Null_Ptr, unless the object is imported (see B.1).

```

function To_Chars_Ptr (Item : in char_array_access; 23
                      Nul_Check : in Boolean := False)
return chars_ptr;

```

If Item is **null**, then To_Chars_Ptr returns Null_Ptr. If Item is not **null**, Nul_Check is True, and Item.all does not contain nul, then the function propagates Terminator_Error; otherwise To_Chars_Ptr performs a pointer conversion with no allocation of memory.

```

function New_Char_Array (Chars : in char_array) return chars_ptr; 25

```

This function returns a pointer to an allocated object initialized to Chars(Chars'First .. Index) & nul, where

- Index = Chars'Last if Chars does not contain nul, or 27
- Index is the smallest size_t value I such that Chars(I+1) = nul. 28

Storage_Error is propagated if the allocation fails.

```

function New_String (Str : in String) return chars_ptr; 29

```

This function is equivalent to New_Char_Array(To_C(Str)).

```

procedure Free (Item : in out chars_ptr); 31

```

If Item is Null_Ptr, then Free has no effect. Otherwise, Free releases the storage occupied by Value(Item), and resets Item to Null_Ptr.

```

function Value (Item : in chars_ptr) return char_array; 33

```

If Item = Null_Ptr then Value propagates Dereference_Error. Otherwise Value returns the prefix of the array of chars pointed to by Item, up to and including the first nul. The lower bound of the result is 0. If Item does not point to a nul-terminated string, then execution of Value is erroneous.

```

function Value (Item : in chars_ptr; Length : in size_t) 35
return char_array;

```

If Item = Null_Ptr then Value propagates Dereference_Error. Otherwise Value returns the shorter of two arrays, either the first Length chars pointed to by Item, or Value(Item). The lower bound of the result is 0. If Length is 0, then Value propagates Constraint_Error.

```

37      function Value (Item : in chars_ptr) return String;
38          Equivalent to To_Ada(Value(Item), Trim_Nul=>True).

39      function Value (Item : in chars_ptr; Length : in size_t)
40/       return String;
41          Equivalent to To_Ada(Value(Item, Length) & nul, Trim_Nul=>True).

42      function Strlen (Item : in chars_ptr) return size_t;
43          Returns Val'Length-1 where Val = Value(Item); propagates Dereference_Error if Item =
44/           Null_Ptr.

43      procedure Update (Item   : in chars_ptr;
44/                      Offset  : in size_t;
45/                      Chars   : in char_array;
46/                      Check   : Boolean := True);
47
48/           If Item = Null_Ptr, then Update propagates Dereference_Error. Otherwise, this procedure
49/           updates the value pointed to by Item, starting at position Offset, using Chars as the data to be
50/           copied into the array. Overwriting the nul terminator, and skipping with the Offset past the nul
51/           terminator, are both prevented if Check is True, as follows:
52/               • Let N = Strlen(Item). If Check is True, then:
53/                   • If Offset+Chars'Length>N, propagate Update_Error.
54/                   • Otherwise, overwrite the data in the array pointed to by Item, starting at the char
55/                       at position Offset, with the data in Chars.
56/               • If Check is False, then processing is as above, but with no check that
57/                 Offset+Chars'Length>N.

59      procedure Update (Item   : in chars_ptr;
60/                      Offset  : in size_t;
61/                      Str     : in String;
62/                      Check   : in Boolean := True);
63
64/           Equivalent to Update(Item, Offset, To_C(Str, Append_Nul=>False), Check).

```

Erroneous Execution

51 Execution of any of the following is erroneous if the Item parameter is not null_ptr and Item does not point to a nul-terminated array of chars.

- 52 • a Value function not taking a Length parameter,
 53 • the Free procedure,
 54 • the Strlen function.

55 Execution of Free(X) is also erroneous if the chars_ptr X was not returned by New_Char_Array or New_String.

56 Reading or updating a freed char_array is erroneous.

57 Execution of Update is erroneous if Check is False and a call with Check equal to True would have propagated Update_Error.

NOTES

58 11 New_Char_Array and New_String might be implemented either through the allocation function from the C environment (“malloc”) or through Ada dynamic memory allocation (“new”). The key points are

- 59 • the returned value (a chars_ptr) is represented as a C “char *” so that it may be passed to C functions;

- the allocated object should be freed by the programmer via a call of Free, not by a called C function.

60

B.3.2 The Generic Package Interfaces.C.Pointers

The generic package Interfaces.C.Pointers allows the Ada programmer to perform C-style operations on pointers. It includes an access type Pointer, Value functions that dereference a Pointer and deliver the designated array, several pointer arithmetic operations, and “copy” procedures that copy the contents of a source pointer into the array designated by a destination pointer. As in C, it treats an object Ptr of type Pointer as a pointer to the first element of an array, so that for example, adding 1 to Ptr yields a pointer to the second element of the array.

The generic allows two styles of usage: one in which the array is terminated by a special terminator element; and another in which the programmer needs to keep track of the length.

Static Semantics

The generic library package Interfaces.C.Pointers has the following declaration:

```

generic
  type Index is (<>);
  type Element is private;
  type Element_Array is array (Index range <>) of aliased Element;
  Default_Terminator : Element;
package Interfaces.C.Pointers is
  pragma Preelaborate(Pointers);

  type Pointer is access all Element;

  function Value(Ref      : in Pointer;
                Terminator : in Element := Default_Terminator)
    return Element_Array;
  function Value(Ref      : in Pointer;
                Length   : in ptrdiff_t)
    return Element_Array;
  Pointer_Error : exception;
  -- C-style Pointer arithmetic
  function "+" (Left : in Pointer; Right : in ptrdiff_t) return Pointer;
  function "+" (Left : in ptrdiff_t; Right : in Pointer)  return Pointer;
  function "-" (Left : in Pointer; Right : in ptrdiff_t) return Pointer;
  function "-" (Left : in Pointer; Right : in Pointer)  return ptrdiff_t;
  procedure Increment (Ref : in out Pointer);
  procedure Decrement (Ref : in out Pointer);
  pragma Convention (Intrinsic, "+");
  pragma Convention (Intrinsic, "-");
  pragma Convention (Intrinsic, Increment);
  pragma Convention (Intrinsic, Decrement);
  function Virtual_Length (Ref      : in Pointer;
                           Terminator : in Element := Default_Terminator)
    return ptrdiff_t;
  procedure Copy_Terminated_Array
    (Source   : in Pointer;
     Target   : in Pointer;
     Limit    : in ptrdiff_t := ptrdiff_t'Last;
     Terminator : in Element := Default_Terminator);
  procedure Copy_Array (Source   : in Pointer;
                       Target   : in Pointer;
                       Length   : in ptrdiff_t);
end Interfaces.C.Pointers;

```

17 The type Pointer is C-compatible and corresponds to one use of C's "Element *". An object of type Pointer
is interpreted as a pointer to the initial Element in an Element_Array. Two styles are supported:

- 18 • Explicit termination of an array value with Default_Terminator (a special terminator value);
- 19 • Programmer-managed length, with Default_Terminator treated simply as a data element.

```
20    function Value(Ref      : in Pointer;
                  Terminator : in Element := Default_Terminator)
                    return Element_Array;
```

21 This function returns an Element_Array whose value is the array pointed to by Ref, up to and
including the first Terminator; the lower bound of the array is IndexFirst.
Interfaces.C.Strings.Dereference_Error is propagated if Ref is **null**.

```
22    function Value(Ref      : in Pointer;
                  Length   : in ptrdiff_t)
                    return Element_Array;
```

23 This function returns an Element_Array comprising the first Length elements pointed to by Ref.
The exception Interfaces.C.Strings.Dereference_Error is propagated if Ref is **null**.

24 The "+" and "--" functions perform arithmetic on Pointer values, based on the Size of the array elements. In
each of these functions, Pointer_Error is propagated if a Pointer parameter is **null**.

```
25    procedure Increment (Ref : in out Pointer);
        Equivalent to Ref := Ref+1.
```

```
27    procedure Decrement (Ref : in out Pointer);
        Equivalent to Ref := Ref-1.
```

```
29    function Virtual_Length (Ref      : in Pointer;
                           Terminator : in Element := Default_Terminator)
                             return ptrdiff_t;
```

30 Returns the number of Elements, up to the one just before the first Terminator, in Value(Ref,
Terminator).

```
31    procedure Copy_Terminated_Array
        (Source      : in Pointer;
         Target     : in Pointer;
         Limit     : in ptrdiff_t := ptrdiff_t'Last;
         Terminator : in Element := Default_Terminator);
```

32 This procedure copies Value(Source, Terminator) into the array pointed to by Target; it stops
either after Terminator has been copied, or the number of elements copied is Limit, whichever
occurs first. Dereference_Error is propagated if either Source or Target is **null**.

```
33    procedure Copy_Array (Source   : in Pointer;
                           Target   : in Pointer;
                           Length   : in ptrdiff_t);
```

34 This procedure copies the first Length elements from the array pointed to by Source, into the
array pointed to by Target. Dereference_Error is propagated if either Source or Target is **null**.

Erroneous Execution

35 It is erroneous to dereference a Pointer that does not designate an aliased Element.

36 Execution of Value(Ref, Terminator) is erroneous if Ref does not designate an aliased Element in an
Element_Array terminated by Terminator.

Execution of Value(Ref, Length) is erroneous if Ref does not designate an aliased Element in an Element_Array containing at least Length Elements between the designated Element and the end of the array, inclusive. 37

Execution of Virtual_Length(Ref, Terminator) is erroneous if Ref does not designate an aliased Element in an Element_Array terminated by Terminator. 38

Execution of Copy_Terminated_Array(Source, Target, Limit, Terminator) is erroneous in either of the following situations: 39

- Execution of both Value(Source, Terminator) and Value(Source, Limit) are erroneous, or 40
- Copying writes past the end of the array containing the Element designated by Target. 41

Execution of Copy_Array(Source, Target, Length) is erroneous if either Value(Source, Length) is erroneous, or copying writes past the end of the array containing the Element designated by Target. 42

NOTES

12 To compose a Pointer from an Element_Array, use 'Access on the first element. For example (assuming appropriate instantiations): 43

```
Some_Array : Element_Array(0..5) ;
Some_Pointer : Pointer := Some_Array(0)'Access;
```

Examples

Example of Interfaces.C.Pointers:

```
with Interfaces.C.Pointers;
with Interfaces.C.Strings;
procedure Test_Pointers is
    package C renames Interfaces.C;
    package Char_Ptrs is
        new C.Pointers (Index          => C.size_t,
                        Element        => C.char,
                        Element_Array   => C.char_array,
                        Default_Terminator => C.nul);
    use type Char_Ptrs.Pointer;
    subtype Char_Star is Char_Ptrs.Pointer;
    procedure Strcpy (Target_Ptr, Source_Ptr : Char_Star) is
        Target_Temp_Ptr : Char_Star := Target_Ptr;
        Source_Temp_Ptr : Char_Star := Source_Ptr;
        Element : C.char;
    begin
        if Target_Temp_Ptr = null or Source_Temp_Ptr = null then
            raise C.Strings.Dereference_Error;
        end if;
        loop
            Element          := Source_Temp_Ptr.all;
            Target_Temp_Ptr.all := Element;
            exit when C."="(Element, C.nul);
            Char_Ptrs.Increment(Target_Temp_Ptr);
            Char_Ptrs.Increment(Source_Temp_Ptr);
        end loop;
    end Strcpy;
begin
    ...
end Test_Pointers;
```

B.3.3 Pragma Unchecked_Union

- 1/2 A pragma `Unchecked_Union` specifies an interface correspondence between a given discriminated type and some C union. The pragma specifies that the associated type shall be given a representation that leaves no space for its discriminant(s).

Syntax

- 2/2 The form of a pragma `Unchecked_Union` is as follows:
- 3/2 **pragma** `Unchecked_Union (first_subtype_local_name);`

Legality Rules

- 4/2 `Unchecked_Union` is a representation pragma, specifying the unchecked union aspect of representation.
- 5/2 The `first_subtype_local_name` of a `pragma` `Unchecked_Union` shall denote an unconstrained discriminated record subtype having a `variant_part`.
- 6/2 A type to which a `pragma` `Unchecked_Union` applies is called an *unchecked union type*. A subtype of an unchecked union type is defined to be an *unchecked union subtype*. An object of an unchecked union type is defined to be an *unchecked union object*.
- 7/2 All component subtypes of an unchecked union type shall be C-compatible.
- 8/2 If a component subtype of an unchecked union type is subject to a per-object constraint, then the component subtype shall be an unchecked union subtype.
- 9/2 Any name that denotes a discriminant of an object of an unchecked union type shall occur within the declarative region of the type.
- 10/2 A component declared in a `variant_part` of an unchecked union type shall not have a controlled, protected, or task part.
- 11/2 The completion of an incomplete or private type declaration having a `known_discriminant_part` shall not be an unchecked union type.
- 12/2 An unchecked union subtype shall only be passed as a generic actual parameter if the corresponding formal type has no known discriminants or is an unchecked union type.

Static Semantics

- 13/2 An unchecked union type is eligible for convention C.
- 14/2 All objects of an unchecked union type have the same size.
- 15/2 Discriminants of objects of an unchecked union type are of size zero.
- 16/2 Any check which would require reading a discriminant of an unchecked union object is suppressed (see 11.5). These checks include:
- 17/2
 - The check performed when addressing a variant component (i.e., a component that was declared in a variant part) of an unchecked union object that the object has this component (see 4.1.3).
 - Any checks associated with a type or subtype conversion of a value of an unchecked union type (see 4.6). This includes, for example, the check associated with the implicit subtype conversion of an assignment statement.
- 18/2

- The subtype membership check associated with the evaluation of a qualified expression (see 4.7) or an uninitialized allocator (see 4.8).

19/2

Dynamic Semantics

A view of an unchecked union object (including a type conversion or function call) has *inferable discriminants* if it has a constrained nominal subtype, unless the object is a component of an enclosing unchecked union object that is subject to a per-object constraint and the enclosing object lacks inferable discriminants.

20/2

An expression of an unchecked union type has inferable discriminants if it is either a name of an object with inferable discriminants or a qualified expression whose `subtype_mark` denotes a constrained subtype.

21/2

`Program_Error` is raised in the following cases:

22/2

- Evaluation of the predefined equality operator for an unchecked union type if either of the operands lacks inferable discriminants.
- Evaluation of the predefined equality operator for a type which has a subcomponent of an unchecked union type whose nominal subtype is unconstrained.
- Evaluation of a membership test if the `subtype_mark` denotes a constrained unchecked union subtype and the expression lacks inferable discriminants.
- Conversion from a derived unchecked union type to an unconstrained non-unchecked-union type if the operand of the conversion lacks inferable discriminants.
- Execution of the default implementation of the `Write` or `Read` attribute of an unchecked union type.
- Execution of the default implementation of the `Output` or `Input` attribute of an unchecked union type if the type lacks default discriminant values.

23/2

24/2

25/2

26/2

27/2

28/2

Implementation Permissions

An implementation may require that `pragma Controlled` be specified for the type of an access subcomponent of an unchecked union type.

29/2

NOTES

13 The use of an unchecked union to obtain the effect of an unchecked conversion results in erroneous execution (see 11.5). Execution of the following example is erroneous even if `Float'Size = Integer'Size`:

30/2

```
type T (Flag : Boolean := False) is
  record
    case Flag is
      when False =>
        F1 : Float := 0.0;
      when True =>
        F2 : Integer := 0;
    end case;
  end record;
pragma Unchecked_Union (T);
X : T;
Y : Integer := X.F2; -- erroneous
```

31/2

32/2

B.4 Interfacing with COBOL

- 1 The facilities relevant to interfacing with the COBOL language are the package Interfaces.COBOL and support for the Import, Export and Convention pragmas with *convention_identifier* COBOL.
- 2 The COBOL interface package supplies several sets of facilities:
- 3 • A set of types corresponding to the native COBOL types of the supported COBOL implementation (so-called “internal COBOL representations”), allowing Ada data to be passed as parameters to COBOL programs
 - 4 • A set of types and constants reflecting external data representations such as might be found in files or databases, allowing COBOL-generated data to be read by an Ada program, and Ada-generated data to be read by COBOL programs
 - 5 • A generic package for converting between an Ada decimal type value and either an internal or external COBOL representation

Static Semantics

- 6 The library package Interfaces.COBOL has the following declaration:

```

7   package Interfaces.COBOL is
8     pragma Preelaborate(COBOL);
9
10    -- Types and operations for internal data representations
11    type Floating      is digits implementation-defined;
12    type Long_Floating is digits implementation-defined;
13
14    type Binary        is range implementation-defined;
15    type Long_Binary   is range implementation-defined;
16
17    Max_Digits_Binary : constant := implementation-defined;
18    Max_Digits_Long_Binary : constant := implementation-defined;
19
20    type Decimal_Element is mod implementation-defined;
21    type Packed_Decimal is array (Positive range <>) of Decimal_Element;
22    pragma Pack(Packed_Decimal);
23
24    type COBOL_Character is implementation-defined character type;
25
26    Ada_To_COBOL : array (Character) of COBOL_Character := implementation-defined;
27    COBOL_To_Ada : array (COBOL_Character) of Character := implementation-defined;
28
29    type Alphanumeric is array (Positive range <>) of COBOL_Character;
30    pragma Pack(Alphanumeric);
31
32    function To_COBOL (Item : in String) return Alphanumeric;
33    function To_Ada    (Item : in Alphanumeric) return String;
34
35    procedure To_COBOL (Item          : in String;
36                         Target       : out Alphanumeric;
37                         Last        : out Natural);
38
39    procedure To_Ada  (Item          : in Alphanumeric;
40                         Target       : out String;
41                         Last        : out Natural);
42
43    type Numeric is array (Positive range <>) of COBOL_Character;
44    pragma Pack(Numeric);
45
46    -- Formats for COBOL data representations
47
48    type Display_Format is private;
49
50    Unsigned           : constant Display_Format;
51    Leading_Separate   : constant Display_Format;
52    Trailing_Separate  : constant Display_Format;
53    Leading_Nonseparate : constant Display_Format;
54    Trailing_Nonseparate : constant Display_Format;

```

```

type Binary_Format is private;                                     24
High_Order_First : constant Binary_Format;                         25
Low_Order_First : constant Binary_Format;
Native_Binary : constant Binary_Format;
type Packed_Format is private;                                     26
Packed_Unsigned : constant Packed_Format;                         27
Packed_Signed : constant Packed_Format;
-- Types for external representation of COBOL binary data           28
type Byte is mod 2**COBOL_Character'Size;                      29
type Byte_Array is array (Positive range <>) of Byte;
pragma Pack (Byte_Array);
Conversion_Error : exception;                                     30
generic
  type Num is delta <> digits <>;
package Decimal_Conversions is
  -- Display Formats: data values are represented as Numeric       31
    function Valid (Item : in Numeric;
                  Format : in Display_Format) return Boolean;          32
    function Length (Format : in Display_Format) return Natural;        33
    function To_Decimal (Item : in Numeric;
                        Format : in Display_Format) return Num;          34
    function To_Display (Item : in Num;
                        Format : in Display_Format) return Numeric;        35
  -- Packed Formats: data values are represented as Packed_Decimal   36
    function Valid (Item : in Packed_Decimal;
                  Format : in Packed_Format) return Boolean;          37
    function Length (Format : in Packed_Format) return Natural;        38
    function To_Decimal (Item : in Packed_Decimal;
                        Format : in Packed_Format) return Num;          39
    function To_Packed (Item : in Num;
                        Format : in Packed_Format) return Packed_Decimal;    40
  -- Binary Formats: external data values are represented as Byte_Array  41
    function Valid (Item : in Byte_Array;
                  Format : in Binary_Format) return Boolean;          42
    function Length (Format : in Binary_Format) return Natural;        43
    function To_Decimal (Item : in Byte_Array;
                        Format : in Binary_Format) return Num;          44
    function To_Binary (Item : in Num;
                        Format : in Binary_Format) return Byte_Array;        45
  -- Internal Binary formats: data values are of type Binary or Long_Binary 46
    function To_Decimal (Item : in Binary) return Num;                47
    function To_Decimal (Item : in Long_Binary) return Num;
    function To_Binary (Item : in Num) return Binary;                 48
    function To_Long_Binary (Item : in Num) return Long_Binary;
  end Decimal_Conversions;                                         49
private
  ... -- not specified by the language                           50
end Interfaces.COBOL;

```

Each of the types in Interfaces.COBOL is COBOL-compatible.

The types Floating and Long_Floating correspond to the native types in COBOL for data items with computational usage implemented by floating point. The types Binary and Long_Binary correspond to the

native types in COBOL for data items with binary usage, or with computational usage implemented by binary.

53 Max_Digits_Binary is the largest number of decimal digits in a numeric value that is represented as Binary. Max_Digits_Long_Binary is the largest number of decimal digits in a numeric value that is represented as Long_Binary.

54 The type Packed.Decimal corresponds to COBOL's packed-decimal usage.

55 The type COBOL_Character defines the run-time character set used in the COBOL implementation. Ada_To_COBOL and COBOL_To_Ada are the mappings between the Ada and COBOL run-time character sets.

56 Type Alphanumeric corresponds to COBOL's alphanumeric data category.

57 Each of the functions To_COBOL and To_Ada converts its parameter based on the mappings Ada_To_COBOL and COBOL_To_Ada, respectively. The length of the result for each is the length of the parameter, and the lower bound of the result is 1. Each component of the result is obtained by applying the relevant mapping to the corresponding component of the parameter.

58 Each of the procedures To_COBOL and To_Ada copies converted elements from Item to Target, using the appropriate mapping (Ada_To_COBOL or COBOL_To_Ada, respectively). The index in Target of the last element assigned is returned in Last (0 if Item is a null array). If Item'Length exceeds Target'Length, Constraint_Error is propagated.

59 Type Numeric corresponds to COBOL's numeric data category with display usage.

60 The types Display_Format, Binary_Format, and Packed_Format are used in conversions between Ada decimal type values and COBOL internal or external data representations. The value of the constant Native_Binary is either High_Order_First or Low_Order_First, depending on the implementation.

```
61      function Valid (Item    : in Numeric;
                      Format   : in Display_Format) return Boolean;
```

62 The function Valid checks that the Item parameter has a value consistent with the value of Format. If the value of Format is other than Unsigned, Leading_Separate, and Trailing_Separate, the effect is implementation defined. If Format does have one of these values, the following rules apply:

- 63/1 • Format=Unsigned: if Item comprises one or more decimal digit characters then Valid returns True, else it returns False.

- 64/1 • Format=Leading_Separate: if Item comprises a single occurrence of the plus or minus sign character, and then one or more decimal digit characters, then Valid returns True, else it returns False.

- 65/1 • Format=Trailing_Separate: if Item comprises one or more decimal digit characters and finally a plus or minus sign character, then Valid returns True, else it returns False.

```
66      function Length (Format : in Display_Format) return Natural;
```

67 The Length function returns the minimal length of a Numeric value sufficient to hold any value of type Num when represented as Format.

```
function To_Decimal (Item : in Numeric;
                     Format : in Display_Format) return Num;
```

68

Produces a value of type Num corresponding to Item as represented by Format. The number of digits after the assumed radix point in Item is Num'Scale. Conversion_Error is propagated if the value represented by Item is outside the range of Num.

69

```
function To_Display (Item : in Num;
                      Format : in Display_Format) return Numeric;
```

70

This function returns the Numeric value for Item, represented in accordance with Format. The length of the returned value is Length(Format), and the lower bound is 1. Conversion_Error is propagated if Num is negative and Format is Unsigned.

71/1

```
function Valid (Item : in Packed_Decimal;
                  Format : in Packed_Format) return Boolean;
```

72

This function returns True if Item has a value consistent with Format, and False otherwise. The rules for the formation of Packed_Decimal values are implementation defined.

73

```
function Length (Format : in Packed_Format) return Natural;
```

74

This function returns the minimal length of a Packed_Decimal value sufficient to hold any value of type Num when represented as Format.

75

```
function To_Decimal (Item : in Packed_Decimal;
                     Format : in Packed_Format) return Num;
```

76

Produces a value of type Num corresponding to Item as represented by Format. Num'Scale is the number of digits after the assumed radix point in Item. Conversion_Error is propagated if the value represented by Item is outside the range of Num.

77

```
function To_Packed (Item : in Num;
                      Format : in Packed_Format) return Packed_Decimal;
```

78

This function returns the Packed_Decimal value for Item, represented in accordance with Format. The length of the returned value is Length(Format), and the lower bound is 1. Conversion_Error is propagated if Num is negative and Format is Packed_Unsigned.

79/1

```
function Valid (Item : in Byte_Array;
                  Format : in Binary_Format) return Boolean;
```

80

This function returns True if Item has a value consistent with Format, and False otherwise.

81

```
function Length (Format : in Binary_Format) return Natural;
```

82

This function returns the minimal length of a Byte_Array value sufficient to hold any value of type Num when represented as Format.

83

```
function To_Decimal (Item : in Byte_Array;
                     Format : in Binary_Format) return Num;
```

84

Produces a value of type Num corresponding to Item as represented by Format. Num'Scale is the number of digits after the assumed radix point in Item. Conversion_Error is propagated if the value represented by Item is outside the range of Num.

85

```
function To_Binary (Item : in Num;
                      Format : in Binary_Format) return Byte_Array;
```

86

This function returns the Byte_Array value for Item, represented in accordance with Format. The length of the returned value is Length(Format), and the lower bound is 1.

87/1

```
88      function To_Decimal (Item : in Binary)      return Num;
89      function To_Decimal (Item : in Long_Binary) return Num;
```

89 These functions convert from COBOL binary format to a corresponding value of the decimal type Num. Conversion_Error is propagated if Item is too large for Num.

```
90      function To_Binary      (Item : in Num)   return Binary;
91      function To_Long_Binary (Item : in Num)   return Long_Binary;
```

91 These functions convert from Ada decimal to COBOL binary format. Conversion_Error is propagated if the value of Item is too large to be represented in the result type.

Implementation Requirements

92 An implementation shall support pragma Convention with a COBOL *convention_identifier* for a COBOL-eligible type (see B.1).

Implementation Permissions

93 An implementation may provide additional constants of the private types Display_Format, Binary_Format, or Packed_Format.

94 An implementation may provide further floating point and integer types in Interfaces.COBOL to match additional native COBOL types, and may also supply corresponding conversion functions in the generic package Decimal_Conversions.

Implementation Advice

95 An Ada implementation should support the following interface correspondences between Ada and COBOL.

- 96 • An Ada **access** T parameter is passed as a “BY REFERENCE” data item of the COBOL type corresponding to T.
- 97 • An Ada **in** scalar parameter is passed as a “BY CONTENT” data item of the corresponding COBOL type.
- 98 • Any other Ada parameter is passed as a “BY REFERENCE” data item of the COBOL type corresponding to the Ada parameter type; for scalars, a local copy is used if necessary to ensure by-copy semantics.

NOTES

99 14 An implementation is not required to support pragma Convention for access types, nor is it required to support pragma Import, Export or Convention for functions.

100 15 If an Ada subprogram is exported to COBOL, then a call from COBOL call may specify either “BY CONTENT” or “BY REFERENCE”.

Examples

101 *Examples of Interfaces.COBOL:*

```
102      with Interfaces.COBOL;
102      procedure Test_Call is
```

```

-- Calling a foreign COBOL program
-- Assume that a COBOL program PROG has the following declaration 103
-- in its LINKAGE section:
-- 01 Parameter-Area
--   05 NAME PIC X(20).
--   05 SSN  PIC X(9).
--   05 SALARY PIC 99999V99 USAGE COMP.
-- The effect of PROG is to update SALARY based on some algorithm

package COBOL renames Interfaces.COBOL; 104
type Salary_Type is delta 0.01 digits 7; 105
type COBOL_Record is 106
  record
    Name   : COBOL.Numeric(1..20);
    SSN    : COBOL.Numeric(1..9);
    Salary : COBOL.Binary; -- Assume Binary = 32 bits
  end record;
pragma Convention (COBOL, COBOL_Record); 107
procedure Prog (Item : in out COBOL_Record);
pragma Import (COBOL, Prog, "PROG");
package Salary_Conversions is 108
  new COBOL.Decimal_Conversions(Salary_Type);
Some_Salary : Salary_Type := 12_345.67; 109
Some_Record : COBOL_Record := 110
  (Name    => "Johnson, John           ",
   SSN     => "111223333",
   Salary  => Salary_Conversions.To_Binary(Some_Salary));
begin
  Prog (Some_Record);
  ...
end Test_Call; 111
with Interfaces.COBOL;
with COBOL.Sequential_IO; -- Assumed to be supplied by implementation
procedure Test_External_Formats is 112
  -- Using data created by a COBOL program
  -- Assume that a COBOL program has created a sequential file with
  -- the following record structure, and that we need to
  -- process the records in an Ada program
  -- 01 EMPLOYEE-RECORD
  --   05 NAME PIC X(20).
  --   05 SSN  PIC X(9).
  --   05 SALARY PIC 99999V99 USAGE COMP.
  --   05 ADJUST PIC S999V999 SIGN LEADING SEPARATE.
  -- The COMP data is binary (32 bits), high-order byte first
  package COBOL renames Interfaces.COBOL; 113
  type Salary_Type      is delta 0.01 digits 7; 114
  type Adjustments_Type is delta 0.001 digits 6;
  type COBOL_Employee_Record_Type is -- External representation 115
    record
      Name   : COBOL.Alphanumeric(1..20);
      SSN    : COBOL.Alphanumeric(1..9);
      Salary : COBOL.Byte_Array(1..4);
      Adjust : COBOL.Numeric(1..7); -- Sign and 6 digits
    end record;
  pragma Convention (COBOL, COBOL_Employee_Record_Type);
  package COBOL_Employee_IO is 116
    new COBOL.Sequential_IO(COBOL_Employee_Record_Type);
  use COBOL_Employee_IO;
  COBOL_File : File_Type; 117

```

```

118      type Ada_Employee_Record_Type is -- Internal representation
119          record
120              Name      : String(1..20);
121              SSN       : String(1..9);
122              Salary    : Salary_Type;
123              Adjust   : Adjustments_Type;
124          end record;
125
126      COBOL_Record : COBOL_Employee_Record_Type;
127      Ada_Record   : Ada_Employee_Record_Type;
128
129      package Salary_Conversions is
130          new COBOL.Decimal_Conversions(Salary_Type);
131      use Salary_Conversions;
132
133      package Adjustments_Conversions is
134          new COBOL.Decimal_Conversions(Adjustments_Type);
135      use Adjustments_Conversions;
136
137      begin
138          Open (COBOL_File, Name => "Some_File");
139
140          loop
141              Read (COBOL_File, COBOL_Record);
142
143              Ada_Record.Name := To_Ada(COBOL_Record.Name);
144              Ada_Record.SSN  := To_Ada(COBOL_Record.SSN);
145              Ada_Record.Salary :=
146                  To_Decimal(COBOL_Record.Salary, COBOL.High_Order_First);
147              Ada_Record.Adjust :=
148                  To_Decimal(COBOL_Record.Adjust, COBOL.Leading_Separate);
149              ... --Process Ada_Record
150          end loop;
151      exception
152          when End_Error => ...
153      end Test_External_Formats;

```

B.5 Interfacing with Fortran

- The facilities relevant to interfacing with the Fortran language are the package Interfaces.Fortran and support for the Import, Export and Convention pragmas with *convention_identifier* Fortran.
- The package Interfaces.Fortran defines Ada types whose representations are identical to the default representations of the Fortran intrinsic types Integer, Real, Double Precision, Complex, Logical, and Character in a supported Fortran implementation. These Ada types can therefore be used to pass objects between Ada and Fortran programs.

Static Semantics

- The library package Interfaces.Fortran has the following declaration:

```

4      with Ada.Numerics.Generic_Complex_Types; -- see G.1.1
5      pragma Elaborate_All(Ada.Numerics.Generic_Complex_Types);
6      package Interfaces.Fortran is
7          pragma Pure(Fortran);
8
9          type Fortran_Integer is range implementation-defined;
10         type Real           is digits implementation-defined;
11         type Double_Precision is digits implementation-defined;
12
13         type Logical is new Boolean;
14
15         package Single_Precision_Complex_Types is
16             new Ada.Numerics.Generic_Complex_Types (Real);
17
18         type Complex is new Single_Precision_Complex_Types.Complex;

```

```

subtype Imaginary is Single_Precision_Complex_Types.Imaginary;          10
i : Imaginary renames Single_Precision_Complex_Types.i;
j : Imaginary renames Single_Precision_Complex_Types.j;
type Character_Set is implementation-defined character type;
type Fortran_Character is array (Positive range <>) of Character_Set;
pragma Pack (Fortran_Character);
function To_Fortran (Item : in Character) return Character_Set;
function To_Ada (Item : in Character_Set) return Character;
function To_Fortran (Item : in String) return Fortran_Character;
function To_Ada (Item : in Fortran_Character) return String;
procedure To_Fortran (Item      : in String;
                      Target    : out Fortran_Character;
                      Last     : out Natural);
procedure To_Ada (Item      : in Fortran_Character;
                      Target    : out String;
                      Last     : out Natural);
end Interfaces.Fortran;                                              17

```

The types Fortran_Integer, Real, Double_Precision, Logical, Complex, and Fortran_Character are Fortran-compatible.

The To_Fortran and To_Ada functions map between the Ada type Character and the Fortran type Character_Set, and also between the Ada type String and the Fortran type Fortran_Character. The To_Fortran and To_Ada procedures have analogous effects to the string conversion subprograms found in Interfaces.COBOL.

Implementation Requirements

An implementation shall support **pragma** Convention with a Fortran *convention_identifier* for a Fortran-eligible type (see B.1).

Implementation Permissions

An implementation may add additional declarations to the Fortran interface packages. For example, the Fortran interface package for an implementation of Fortran 77 (ANSI X3.9-1978) that defines types like Integer*n, Real*n, Logical*n, and Complex*n may contain the declarations of types named Integer_Star_n, Real_Star_n, Logical_Star_n, and Complex_Star_n. (This convention should not apply to Character*n, for which the Ada analog is the constrained array subtype Fortran_Character (1..n).) Similarly, the Fortran interface package for an implementation of Fortran 90 that provides multiple *kinds* of intrinsic types, e.g. Integer (Kind=n), Real (Kind=n), Logical (Kind=n), Complex (Kind=n), and Character (Kind=n), may contain the declarations of types with the recommended names Integer_Kind_n, Real_Kind_n, Logical_Kind_n, Complex_Kind_n, and Character_Kind_n.

Implementation Advice

An Ada implementation should support the following interface correspondences between Ada and Fortran:

- An Ada procedure corresponds to a Fortran subroutine.
- An Ada function corresponds to a Fortran function.
- An Ada parameter of an elementary, array, or record type T is passed as a T_F argument to a Fortran procedure, where T_F is the Fortran type corresponding to the Ada type T, and where the INTENT attribute of the corresponding dummy argument matches the Ada formal parameter mode; the Fortran implementation's parameter passing conventions are used. For elementary types, a local copy is used if necessary to ensure by-copy semantics.

- 26 • An Ada parameter of an access-to-subprogram type is passed as a reference to a Fortran procedure whose interface corresponds to the designated subprogram's specification.

NOTES

- 27 16 An object of a Fortran-compatible record type, declared in a library package or subprogram, can correspond to a Fortran common block; the type also corresponds to a Fortran “derived type”.

Examples

28 Example of *Interfaces.Fortran*:

```

29   with Interfaces.Fortran;
30   use Interfaces.Fortran;
31   procedure Ada_Application is
32     type Fortran_Matrix is array (Integer range <>,
33                                 Integer range <>) of Double_Precision;
34     pragma Convention (Fortran, Fortran_Matrix);      -- stored in Fortran's
35                                         -- column-major order
36     procedure Invert (Rank : in Fortran_Integer; X : in out Fortran_Matrix);
37     pragma Import (Fortran, Invert);                  -- a Fortran subroutine
38     Rank      : constant Fortran_Integer := 100;
39     My_Matrix : Fortran_Matrix (1 .. Rank, 1 .. Rank);
40
41 begin
42   ...
43   My_Matrix := ...;
44   ...
45   Invert (Rank, My_Matrix);
46   ...
47 end Ada_Application;
```

Annex C (normative) Systems Programming

The Systems Programming Annex specifies additional capabilities provided for low-level programming. These capabilities are also required in many real-time, embedded, distributed, and information systems.

C.1 Access to Machine Operations

This clause specifies rules regarding access to machine instructions from within an Ada program.

1

Implementation Requirements

The implementation shall support machine code insertions (see 13.8) or intrinsic subprograms (see 6.3.1) (or both). Implementation-defined attributes shall be provided to allow the use of Ada entities as operands.

2

Implementation Advice

The machine code or intrinsics support should allow access to all operations normally available to assembly language programmers for the target environment, including privileged instructions, if any.

3

The interfacing pragmas (see Annex B) should support interface to assembler; the default assembler should be associated with the convention identifier Assembler.

4

If an entity is exported to assembly language, then the implementation should allocate it at an addressable location, and should ensure that it is retained by the linking process, even if not otherwise referenced from the Ada code. The implementation should assume that any call to a machine code or assembler subprogram is allowed to read or update every object that is specified as exported.

5

Documentation Requirements

The implementation shall document the overhead associated with calling machine-code or intrinsic subprograms, as compared to a fully-inlined call, and to a regular out-of-line call.

6

The implementation shall document the types of the package System.Machine_Code usable for machine code insertions, and the attributes to be used in machine code insertions for references to Ada entities.

7

The implementation shall document the subprogram calling conventions associated with the convention identifiers available for use with the interfacing pragmas (Ada and Assembler, at a minimum), including register saving, exception propagation, parameter passing, and function value returning.

8

For exported and imported subprograms, the implementation shall document the mapping between the Link_Name string, if specified, or the Ada designator, if not, and the external link name used for such a subprogram.

9

Implementation Advice

The implementation should ensure that little or no overhead is associated with calling intrinsic and machine-code subprograms.

10

It is recommended that intrinsic subprograms be provided for convenient access to any machine operations that provide special capabilities or efficiency and that are not otherwise available through the language constructs. Examples of such instructions include:

11

- 12 • Atomic read-modify-write operations — e.g., test and set, compare and swap, decrement and
test, enqueue/dequeue.
- 13 • Standard numeric functions — e.g., *sin*, *log*.
- 14 • String manipulation operations — e.g., translate and test.
- 15 • Vector operations — e.g., compare vector against thresholds.
- 16 • Direct operations on I/O ports.

C.2 Required Representation Support

1/2 This clause specifies minimal requirements on the support for representation items and related features.

Implementation Requirements

- 2 The implementation shall support at least the functionality defined by the recommended levels of support in Section 13.

C.3 Interrupt Support

1 This clause specifies the language-defined model for hardware interrupts in addition to mechanisms for handling interrupts.

Dynamic Semantics

- 2 An *interrupt* represents a class of events that are detected by the hardware or the system software. Interrupts are said to occur. An *occurrence* of an interrupt is separable into generation and delivery. *Generation* of an interrupt is the event in the underlying hardware or system that makes the interrupt available to the program. *Delivery* is the action that invokes part of the program as response to the interrupt occurrence. Between generation and delivery, the interrupt occurrence (or interrupt) is *pending*. Some or all interrupts may be *blocked*. When an interrupt is blocked, all occurrences of that interrupt are prevented from being delivered. Certain interrupts are *reserved*. The set of reserved interrupts is implementation defined. A reserved interrupt is either an interrupt for which user-defined handlers are not supported, or one which already has an attached handler by some other implementation-defined means. Program units can be connected to non-reserved interrupts. While connected, the program unit is said to be *attached* to that interrupt. The execution of that program unit, the *interrupt handler*, is invoked upon delivery of the interrupt occurrence.

3 While a handler is attached to an interrupt, it is called once for each delivered occurrence of that interrupt. While the handler executes, the corresponding interrupt is blocked.

4 While an interrupt is blocked, all occurrences of that interrupt are prevented from being delivered. Whether such occurrences remain pending or are lost is implementation defined.

5 Each interrupt has a *default treatment* which determines the system's response to an occurrence of that interrupt when no user-defined handler is attached. The set of possible default treatments is implementation defined, as is the method (if one exists) for configuring the default treatments for interrupts.

6 An interrupt is delivered to the handler (or default treatment) that is in effect for that interrupt at the time of delivery.

7 An exception propagated from a handler that is invoked by an interrupt has no effect.

If the Ceiling_Locking policy (see D.3) is in effect, the interrupt handler executes with the active priority that is the ceiling priority of the corresponding protected object. 8

Implementation Requirements

The implementation shall provide a mechanism to determine the minimum stack space that is needed for each interrupt handler and to reserve that space for the execution of the handler. This space should accommodate nested invocations of the handler where the system permits this. 9

If the hardware or the underlying system holds pending interrupt occurrences, the implementation shall provide for later delivery of these occurrences to the program. 10

If the Ceiling_Locking policy is not in effect, the implementation shall provide means for the application to specify whether interrupts are to be blocked during protected actions. 11

Documentation Requirements

The implementation shall document the following items: 12

1. For each interrupt, which interrupts are blocked from delivery when a handler attached to that interrupt executes (either as a result of an interrupt delivery or of an ordinary call on a procedure of the corresponding protected object). 13
2. Any interrupts that cannot be blocked, and the effect of attaching handlers to such interrupts, if this is permitted. 14
3. Which run-time stack an interrupt handler uses when it executes as a result of an interrupt delivery; if this is configurable, what is the mechanism to do so; how to specify how much space to reserve on that stack. 15
4. Any implementation- or hardware-specific activity that happens before a user-defined interrupt handler gets control (e.g., reading device registers, acknowledging devices). 16
5. Any timing or other limitations imposed on the execution of interrupt handlers. 17
6. The state (blocked/unblocked) of the non-reserved interrupts when the program starts; if some interrupts are unblocked, what is the mechanism a program can use to protect itself before it can attach the corresponding handlers. 18
7. Whether the interrupted task is allowed to resume execution before the interrupt handler returns. 19
8. The treatment of interrupt occurrences that are generated while the interrupt is blocked; i.e., whether one or more occurrences are held for later delivery, or all are lost. 20
9. Whether predefined or implementation-defined exceptions are raised as a result of the occurrence of any interrupt, and the mapping between the machine interrupts (or traps) and the predefined exceptions. 21
10. On a multi-processor, the rules governing the delivery of an interrupt to a particular processor. 22

Implementation Permissions

If the underlying system or hardware does not allow interrupts to be blocked, then no blocking is required as part of the execution of subprograms of a protected object for which one of its subprograms is an interrupt handler. 23/2

In a multi-processor with more than one interrupt subsystem, it is implementation defined whether (and how) interrupt sources from separate subsystems share the same Interrupt_ID type (see C.3.2). In particular, the meaning of a blocked or pending interrupt may then be applicable to one processor only. 24

- 25 Implementations are allowed to impose timing or other limitations on the execution of interrupt handlers.
- 26/2 Other forms of handlers are allowed to be supported, in which case the rules of this clause should be adhered to.
- 27 The active priority of the execution of an interrupt handler is allowed to vary from one occurrence of the same interrupt to another.

Implementation Advice

- 28/2 If the Ceiling_Locking policy is not in effect, the implementation should provide means for the application to specify which interrupts are to be blocked during protected actions, if the underlying system allows for finer-grained control of interrupt blocking.

NOTES

1 The default treatment for an interrupt can be to keep the interrupt pending or to deliver it to an implementation-defined handler. Examples of actions that an implementation-defined handler is allowed to perform include aborting the partition, ignoring (i.e., discarding occurrences of) the interrupt, or queuing one or more occurrences of the interrupt for possible later delivery when a user-defined handler is attached to that interrupt.

2 It is a bounded error to call Task_Identification.Current_Task (see C.7.1) from an interrupt handler.

3 The rule that an exception propagated from an interrupt handler has no effect is modeled after the rule about exceptions propagated out of task bodies.

C.3.1 Protected Procedure Handlers

Syntax

1 The form of a **pragma** Interrupt_Handler is as follows:

pragma Interrupt_Handler(*handler_name*);

2 The form of a **pragma** Attach_Handler is as follows:

pragma Attach_Handler(*handler_name*, *expression*);

Name Resolution Rules

5 For the Interrupt_Handler and Attach_Handler pragmas, the *handler_name* shall resolve to denote a protected procedure with a parameterless profile.

6 For the Attach_Handler pragma, the expected type for the *expression* is Interrupts.Interrupt_ID (see C.3.2).

Legality Rules

7/2 The Attach_Handler pragma is only allowed immediately within the **protected_definition** where the corresponding subprogram is declared. The corresponding **protected_type_declaration** or **single_protected_declaration** shall be a library-level declaration.

8/2 The Interrupt_Handler pragma is only allowed immediately within the **protected_definition** where the corresponding subprogram is declared. The corresponding **protected_type_declaration** or **single_protected_declaration** shall be a library-level declaration.

Dynamic Semantics

9 If the **pragma** Interrupt_Handler appears in a **protected_definition**, then the corresponding procedure can be attached dynamically, as a handler, to interrupts (see C.3.2). Such procedures are allowed to be attached to multiple interrupts.

The expression in the `Attach_Handler` pragma as evaluated at object creation time specifies an interrupt. As part of the initialization of that object, if the `Attach_Handler` pragma is specified, the `handler` procedure is attached to the specified interrupt. A check is made that the corresponding interrupt is not reserved. `Program_Error` is raised if the check fails, and the existing treatment for the interrupt is not affected.

If the `Ceiling_Locking` policy (see D.3) is in effect, then upon the initialization of a protected object for which either an `Attach_Handler` or `Interrupt_Handler` pragma applies to one of its procedures, a check is made that the ceiling priority defined in the `protected_definition` is in the range of `System Interrupt_Priority`. If the check fails, `Program_Error` is raised.

When a protected object is finalized, for any of its procedures that are attached to interrupts, the handler is detached. If the handler was attached by a procedure in the `Interrupts` package or if no user handler was previously attached to the interrupt, the default treatment is restored. If an `Attach_Handler` pragma was used and the most recently attached handler for the same interrupt is the same as the one that was attached at the time the protected object was initialized, the previous handler is restored.

When a handler is attached to an interrupt, the interrupt is blocked (subject to the Implementation Permission in C.3) during the execution of every protected action on the protected object containing the handler.

Erroneous Execution

If the `Ceiling_Locking` policy (see D.3) is in effect and an interrupt is delivered to a handler, and the interrupt hardware priority is higher than the ceiling priority of the corresponding protected object, the execution of the program is erroneous.

If the handlers for a given interrupt attached via pragma `Attach_Handler` are not attached and detached in a stack-like (LIFO) order, program execution is erroneous. In particular, when a protected object is finalized, the execution is erroneous if any of the procedures of the protected object are attached to interrupts via pragma `Attach_Handler` and the most recently attached handler for the same interrupt is not the same as the one that was attached at the time the protected object was initialized.

Metrics

The following metric shall be documented by the implementation:

- The worst-case overhead for an interrupt handler that is a parameterless protected procedure, in clock cycles. This is the execution time not directly attributable to the handler procedure or the interrupted execution. It is estimated as $C - (A+B)$, where A is how long it takes to complete a given sequence of instructions without any interrupt, B is how long it takes to complete a normal call to a given protected procedure, and C is how long it takes to complete the same sequence of instructions when it is interrupted by one execution of the same procedure called via an interrupt.

Implementation Permissions

When the pragmas `Attach_Handler` or `Interrupt_Handler` apply to a protected procedure, the implementation is allowed to impose implementation-defined restrictions on the corresponding `protected_type_declaration` and `protected_body`.

An implementation may use a different mechanism for invoking a protected procedure in response to a hardware interrupt than is used for a call to that protected procedure from a task.

Notwithstanding what this subclause says elsewhere, the `Attach_Handler` and `Interrupt_Handler` pragmas are allowed to be used for other, implementation defined, forms of interrupt handlers.

Implementation Advice

- 20 Whenever possible, the implementation should allow interrupt handlers to be called directly by the hardware.
- 21 Whenever practical, the implementation should detect violations of any implementation-defined restrictions before run time.

NOTES

- 22 4 The `Attach_Handler` pragma can provide static attachment of handlers to interrupts if the implementation supports preelaboration of protected objects. (See C.4.)
- 23/2 5 A protected object that has a (protected) procedure attached to an interrupt should have a ceiling priority at least as high as the highest processor priority at which that interrupt will ever be delivered.
- 24 6 Protected procedures can also be attached dynamically to interrupts via operations declared in the predefined package `Interrupts`.
- 25 7 An example of a possible implementation-defined restriction is disallowing the use of the standard storage pools within the body of a protected procedure that is an interrupt handler.

C.3.2 The Package Interrupts*Static Semantics*

- 1 The following language-defined packages exist:

```

2   with System;
3   package Ada.Interrupts is
4     type Interrupt_ID is implementation-defined;
5     type Parameterless_Handler is
6       access protected procedure;
7
8     This paragraph was deleted.
9     function Is_Reserved (Interrupt : Interrupt_ID)
10    return Boolean;
11
12     function Is_Attached (Interrupt : Interrupt_ID)
13    return Boolean;
14
15     function Current_Handler (Interrupt : Interrupt_ID)
16    return Parameterless_Handler;
17
18     procedure Attach_Handler
19       (New_Handler : in Parameterless_Handler;
20        Interrupt : in Interrupt_ID);
21
22     procedure Exchange_Handler
23       (Old_Handler : out Parameterless_Handler;
24        New_Handler : in Parameterless_Handler;
25        Interrupt : in Interrupt_ID);
26
27     procedure Detach_Handler
28       (Interrupt : in Interrupt_ID);
29
30     function Reference(Interrupt : Interrupt_ID)
31    return System.Address;
32
33     private
34       ... -- not specified by the language
35   end Ada.Interrupts;
36
37   package Ada.Interrupts.Names is
38     implementation-defined : constant Interrupt_ID := 
39       implementation-defined;
40
41     implementation-defined : constant Interrupt_ID := 
42       implementation-defined;
43   end Ada.Interrupts.Names;

```

Dynamic Semantics

The Interrupt_ID type is an implementation-defined discrete type used to identify interrupts.	13
The Is_Reserved function returns True if and only if the specified interrupt is reserved.	14
The Is_Attached function returns True if and only if a user-specified interrupt handler is attached to the interrupt.	15
The Current_Handler function returns a value that represents the attached handler of the interrupt. If no user-defined handler is attached to the interrupt, Current_Handler returns null .	16/1
The Attach_Handler procedure attaches the specified handler to the interrupt, overriding any existing treatment (including a user handler) in effect for that interrupt. If New_Handler is null , the default treatment is restored. If New_Handler designates a protected procedure to which the pragma Interrupt_Handler does not apply, Program_Error is raised. In this case, the operation does not modify the existing interrupt treatment.	17
The Exchange_Handler procedure operates in the same manner as Attach_Handler with the addition that the value returned in Old_Handler designates the previous treatment for the specified interrupt. If the previous treatment is not a user-defined handler, null is returned.	18/1
The Detach_Handler procedure restores the default treatment for the specified interrupt.	19
For all operations defined in this package that take a parameter of type Interrupt_ID, with the exception of Is_Reserved and Reference, a check is made that the specified interrupt is not reserved. Program_Error is raised if this check fails.	20
If, by using the Attach_Handler, Detach_Handler, or Exchange_Handler procedures, an attempt is made to detach a handler that was attached statically (using the pragma Attach_Handler), the handler is not detached and Program_Error is raised.	21
The Reference function returns a value of type System.Address that can be used to attach a task entry via an address clause (see J.7.1) to the interrupt specified by Interrupt. This function raises Program_Error if attaching task entries to interrupts (or to this particular interrupt) is not supported.	22/2

Implementation Requirements

At no time during attachment or exchange of handlers shall the current handler of the corresponding interrupt be undefined.	23
---	----

Documentation Requirements

If the Ceiling_Locking policy (see D.3) is in effect, the implementation shall document the default ceiling priority assigned to a protected object that contains either the Attach_Handler or Interrupt_Handler pragmas, but not the Interrupt_Priority pragma. This default need not be the same for all interrupts.	24/2
--	------

Implementation Advice

If implementation-defined forms of interrupt handler procedures are supported, such as protected procedures with parameters, then for each such form of a handler, a type analogous to Parameterless_Handler should be specified in a child package of Interrupts, with the same operations as in the predefined package Interrupts.	25
--	----

NOTES

8 The package Interrupts.Names contains implementation-defined names (and constant values) for the interrupts that are supported by the implementation.	26
---	----

Examples

27 Example of interrupt handlers:

```

28      Device_Priority : constant
          array (1..5) of System Interrupt_Priority := ( ... );
protected type Device_Interface
  (Int_ID : Ada Interrupts Interrupt_ID) is
  procedure Handler;
  pragma Attach_Handler(Handler, Int_ID);
  ...
  pragma Interrupt_Priority(Device_Priority(Int_ID));
end Device_Interface;
...
Device_1_Driver : Device_Interface(1);
...
Device_5_Driver : Device_Interface(5);
...

```

C.4 Preelaboration Requirements

1 This clause specifies additional implementation and documentation requirements for the Preelaborate pragma (see 10.2.1).

Implementation Requirements

- 2 The implementation shall not incur any run-time overhead for the elaboration checks of subprograms and **protected bodies** declared in preelaborated library units.
- 3 The implementation shall not execute any memory write operations after load time for the elaboration of constant objects declared immediately within the declarative region of a preelaborated library package, so long as the subtype and initial expression (or default initial expressions if initialized by default) of the **object_declaration** satisfy the following restrictions. The meaning of *load time* is implementation defined.
 - 4
 - Any **subtype_mark** denotes a statically constrained subtype, with statically constrained subcomponents, if any;
 - 4.1/2 • no **subtype_mark** denotes a controlled type, a private type, a private extension, a generic formal private type, a generic formal derived type, or a descendant of such a type;
 - 5 • any **constraint** is a static constraint;
 - 6 • any **allocator** is for an access-to-constant type;
 - 7 • any uses of predefined operators appear only within static expressions;
 - 8 • any **primaries** that are **names**, other than **attribute_references** for the Access or Address attributes, appear only within static expressions;
 - 9 • any **name** that is not part of a static expression is an expanded name or **direct_name** that statically denotes some entity;
 - 10 • any **discrete_choice** of an **array_aggregate** is static;
 - 11 • no language-defined check associated with the elaboration of the **object_declaration** can fail.

Documentation Requirements

- 12 The implementation shall document any circumstances under which the elaboration of a preelaborated package causes code to be executed at run time.
- 13 The implementation shall document whether the method used for initialization of preelaborated variables allows a partition to be restarted without reloading.

Implementation Advice

It is recommended that preelaborated packages be implemented in such a way that there should be little or no code executed at run time for the elaboration of entities not already covered by the Implementation Requirements.

14

C.5 Pragma Discard_Names

A **pragma Discard_Names** may be used to request a reduction in storage used for the names of certain entities.

1

Syntax

The form of a **pragma Discard_Names** is as follows:

2

```
pragma Discard_Names[([On => ] local_name)];
```

3

A **pragma Discard_Names** is allowed only immediately within a **declarative_part**, immediately within a **package_specification**, or as a configuration pragma.

4

Legality Rules

The **local_name** (if present) shall denote a non-derived enumeration first subtype, a tagged first subtype, or an exception. The pragma applies to the type or exception. Without a **local_name**, the pragma applies to all such entities declared after the pragma, within the same declarative region. Alternatively, the pragma can be used as a configuration pragma. If the pragma applies to a type, then it applies also to all descendants of the type.

5

Static Semantics

If a **local_name** is given, then a **pragma Discard_Names** is a representation pragma.

6

If the pragma applies to an enumeration type, then the semantics of the **Wide_Wide_Image** and **Wide_Wide_Value** attributes are implementation defined for that type; the semantics of **Image**, **Wide_Image**, **Value**, and **Wide_Value** are still defined in terms of **Wide_Wide_Image** and **Wide_Wide_Value**. In addition, the semantics of **Text_IO.Enumeration_IO** are implementation defined. If the pragma applies to a tagged type, then the semantics of the **Tags.Wide_Wide_Expanded_Name** function are implementation defined for that type; the semantics of **Tags.Expanded_Name** and **Tags.Wide_Expanded_Name** are still defined in terms of **Tags.Wide_Wide_Expanded_Name**. If the pragma applies to an exception, then the semantics of the **Exceptions.Wide_Wide_Exception_Name** function are implementation defined for that exception; the semantics of **Exceptions.Exception_Name** and **Exceptions.Wide_Exception_Name** are still defined in terms of **Exceptions.Wide_Wide_Exception_Name**.

7/2

Implementation Advice

If the pragma applies to an entity, then the implementation should reduce the amount of storage used for storing names associated with that entity.

8

C.6 Shared Variable Control

This clause specifies representation pragmas that control the use of shared variables.

1

Syntax

The form for pragmas **Atomic**, **Volatile**, **Atomic_Components**, and **Volatile_Components** is as follows:

2

```

3      pragma Atomic(local_name);
4      pragma Volatile(local_name);
5      pragma Atomic_Components(array_local_name);
6      pragma Volatile_Components(array_local_name);

```

- 7/2 An *atomic* type is one to which a pragma Atomic applies. An *atomic* object (including a component) is one to which a pragma Atomic applies, or a component of an array to which a pragma Atomic_Components applies, or any object of an atomic type, other than objects obtained by evaluating a slice.
- 8 A *volatile* type is one to which a pragma Volatile applies. A *volatile* object (including a component) is one to which a pragma Volatile applies, or a component of an array to which a pragma Volatile_Components applies, or any object of a volatile type. In addition, every atomic type or object is also defined to be volatile. Finally, if an object is volatile, then so are all of its subcomponents (the same does not apply to atomic).

Name Resolution Rules

- 9 The local_name in an Atomic or Volatile pragma shall resolve to denote either an object_declaration, a non-inherited component_declaration, or a full_type_declaration. The array_local_name in an Atomic_Components or Volatile_Components pragma shall resolve to denote the declaration of an array type or an array object of an anonymous type.

Legality Rules

- 10 It is illegal to apply either an Atomic or Atomic_Components pragma to an object or type if the implementation cannot support the indivisible reads and updates required by the pragma (see below).
- 11 It is illegal to specify the Size attribute of an atomic object, the Component_Size attribute for an array type with atomic components, or the layout attributes of an atomic component, in a way that prevents the implementation from performing the required indivisible reads and updates.
- 12 If an atomic object is passed as a parameter, then the type of the formal parameter shall either be atomic or allow pass by copy (that is, not be a nonatomic by-reference type). If an atomic object is used as an actual for a generic formal object of mode **in out**, then the type of the generic formal object shall be atomic. If the prefix of an attribute_reference for an Access attribute denotes an atomic object (including a component), then the designated type of the resulting access type shall be atomic. If an atomic type is used as an actual for a generic formal derived type, then the ancestor of the formal type shall be atomic or allow pass by copy. Corresponding rules apply to volatile objects and types.
- 13 If a pragma Volatile, Volatile_Components, Atomic, or Atomic_Components applies to a stand-alone constant object, then a pragma Import shall also apply to it.

Static Semantics

- 14 These pragmas are representation pragmas (see 13.1).

Dynamic Semantics

- 15 For an atomic object (including an atomic component) all reads and updates of the object as a whole are indivisible.
- 16 For a volatile object all reads and updates of the object as a whole are performed directly to memory.
- 17 Two actions are sequential (see 9.10) if each is the read or update of the same atomic object.

If a type is atomic or volatile and it is not a by-copy type, then the type is defined to be a by-reference type. If any subcomponent of a type is atomic or volatile, then the type is defined to be a by-reference type.

If an actual parameter is atomic or volatile, and the corresponding formal parameter is not, then the parameter is passed by copy.

Implementation Requirements

The external effect of a program (see 1.1.3) is defined to include each read and update of a volatile or atomic object. The implementation shall not generate any memory reads or updates of atomic or volatile objects other than those specified by the program.

If a pragma Pack applies to a type any of whose subcomponents are atomic, the implementation shall not pack the atomic subcomponents more tightly than that for which it can support indivisible reads and updates.

Implementation Advice

A load or store of a volatile object whose size is a multiple of System.Storage_Unit and whose alignment is nonzero, should be implemented by accessing exactly the bits of the object and no others.

A load or store of an atomic object should, where possible, be implemented by a single load or store instruction.

NOTES

9 An imported volatile or atomic constant behaves as a constant (i.e. read-only) with respect to other parts of the Ada program, but can still be modified by an “external source.”

C.7 Task Information

This clause describes operations and attributes that can be used to obtain the identity of a task. In addition, a package that associates user-defined information with a task is defined. Finally, a package that associates termination procedures with a task or set of tasks is defined.

C.7.1 The Package Task_Identification

Static Semantics

The following language-defined library package exists:

```
package Ada.Task_Identification is
  pragma Preelaborate(Task_Identification);
  type Task_Id is private;
  pragma Preelaborable_Initialization (Task_Id);
  Null_Task_Id : constant Task_Id;
  function "=" (Left, Right : Task_Id) return Boolean;
  function Image      (T : Task_Id) return String;
  function Current_Task return Task_Id;
  procedure Abort_Task (T : in Task_Id);
  function Is_Terminated(T : Task_Id) return Boolean;
  function Is_Callable  (T : Task_Id) return Boolean;
private
  ... -- not specified by the language
end Ada.Task_Identification;
```

Dynamic Semantics

- 5 A value of the type Task_Id identifies an existent task. The constant Null_Task_Id does not identify any task. Each object of the type Task_Id is default initialized to the value of Null_Task_Id.
- 6 The function "=" returns True if and only if Left and Right identify the same task or both have the value Null_Task_Id.
- 7 The function Image returns an implementation-defined string that identifies T. If T equals Null_Task_Id, Image returns an empty string.
- 8 The function Current_Task returns a value that identifies the calling task.
- 9 The effect of Abort_Task is the same as the `abort_statement` for the task identified by T. In addition, if T identifies the environment task, the entire partition is aborted. See E.1.
- 10 The functions Is_Terminated and Is_Callable return the value of the corresponding attribute of the task identified by T.
- 11 For a prefix T that is of a task type (after any implicit dereference), the following attribute is defined:
 - 12 T'Identity Yields a value of the type Task_Id that identifies the task denoted by T.
- 13 For a prefix E that denotes an `entry_declaration`, the following attribute is defined:
 - 14 E'Caller Yields a value of the type Task_Id that identifies the task whose call is now being serviced. Use of this attribute is allowed only inside an `entry_body` or `accept_statement` corresponding to the `entry_declaration` denoted by E.
- 15 Program_Error is raised if a value of Null_Task_Id is passed as a parameter to Abort_Task, Is_Terminated, and Is_Callable.
- 16 Abort_Task is a potentially blocking operation (see 9.5.1).

Bounded (Run-Time) Errors

- 17/2 It is a bounded error to call the Current_Task function from an entry body, interrupt handler, or finalization of a task attribute. Program_Error is raised, or an implementation-defined value of the type Task_Id is returned.

Erroneous Execution

- 18 If a value of Task_Id is passed as a parameter to any of the operations declared in this package (or any language-defined child of this package), and the corresponding task object no longer exists, the execution of the program is erroneous.

Documentation Requirements

- 19 The implementation shall document the effect of calling Current_Task from an entry body or interrupt handler.

NOTES

- 20 10 This package is intended for use in writing user-defined task scheduling packages and constructing server tasks. Current_Task can be used in conjunction with other operations requiring a task as an argument such as Set_Priority (see D.5).
- 21 11 The function Current_Task and the attribute Caller can return a Task_Id value that identifies the environment task.

C.7.2 The Package Task_Attributes

Static Semantics

The following language-defined generic library package exists:

```

1   with Ada.Task_Identification; use Ada.Task_Identification;
2   generic
3     type Attribute is private;
4     Initial_Value : in Attribute;
5   package Ada.Task_Attributes is
6     type Attribute_Handle is access all Attribute;
7     function Value(T : Task_Id := Current_Task)
8       return Attribute;
9     function Reference(T : Task_Id := Current_Task)
10      return Attribute_Handle;
11    procedure Set_Value(Val : in Attribute;
12                          T : in Task_Id := Current_Task);
13    procedure Reinitialize(T : in Task_Id := Current_Task);
14  end Ada.Task_Attributes;
15

```

Dynamic Semantics

When an instance of Task_Attributes is elaborated in a given active partition, an object of the actual type corresponding to the formal type Attribute is implicitly created for each task (of that partition) that exists and is not yet terminated. This object acts as a user-defined attribute of the task. A task created previously in the partition and not yet terminated has this attribute from that point on. Each task subsequently created in the partition will have this attribute when created. In all these cases, the initial value of the given attribute is Initial_Value.

The Value operation returns the value of the corresponding attribute of T.

The Reference operation returns an access value that designates the corresponding attribute of T.

The Set_Value operation performs any finalization on the old value of the attribute of T and assigns Val to that attribute (see 5.2 and 7.6).

The effect of the Reinitialize operation is the same as Set_Value where the Val parameter is replaced with Initial_Value.

For all the operations declared in this package, Tasking_Error is raised if the task identified by T is terminated. Program_Error is raised if the value of T is Null_Task_Id.

After a task has terminated, all of its attributes are finalized, unless they have been finalized earlier. When the master of an instantiation of Ada.Task_Attributes is finalized, the corresponding attribute of each task is finalized, unless it has been finalized earlier.

Bounded (Run-Time) Errors

If the package Ada.Task_Attributes is instantiated with a controlled type and the controlled type has user-defined Adjust or Finalize operations that in turn access task attributes by any of the above operations, then a call of Set_Value of the instantiated package constitutes a bounded error. The call may perform as expected or may result in forever blocking the calling task and subsequently some or all tasks of the partition.

Erroneous Execution

- 14 It is erroneous to dereference the access value returned by a given call on Reference after a subsequent call on Reinitialize for the same task attribute, or after the associated task terminates.
- 15 If a value of Task_Id is passed as a parameter to any of the operations declared in this package and the corresponding task object no longer exists, the execution of the program is erroneous.
- 15.1/2 An access to a task attribute via a value of type Attribute_Handle is erroneous if executed concurrently with another such access or a call of any of the operations declared in package Task_Attributes. An access to a task attribute is erroneous if executed concurrently with or after the finalization of the task attribute.

Implementation Requirements

- 16/1 For a given attribute of a given task, the implementation shall perform the operations declared in this package atomically with respect to any of these operations of the same attribute of the same task. The granularity of any locking mechanism necessary to achieve such atomicity is implementation defined.
- 17/2 After task attributes are finalized, the implementation shall reclaim any storage associated with the attributes.

Documentation Requirements

- 18 The implementation shall document the limit on the number of attributes per task, if any, and the limit on the total storage for attribute values per task, if such a limit exists.
- 19 In addition, if these limits can be configured, the implementation shall document how to configure them.

Metrics

- 20/2 The implementation shall document the following metrics: A task calling the following subprograms shall execute at a sufficiently high priority as to not be preempted during the measurement period. This period shall start just before issuing the call and end just after the call completes. If the attributes of task T are accessed by the measurement tests, no other task shall access attributes of that task during the measurement period. For all measurements described here, the Attribute type shall be a scalar type whose size is equal to the size of the predefined type Integer. For each measurement, two cases shall be documented: one where the accessed attributes are of the calling task (that is, the default value for the T parameter is used), and the other, where T identifies another, non-terminated, task.
- 21 The following calls (to subprograms in the Task_Attributes package) shall be measured:
 - 22 • a call to Value, where the return value is Initial_Value;
 - 23 • a call to Value, where the return value is not equal to Initial_Value;
 - 24 • a call to Reference, where the return value designates a value equal to Initial_Value;
 - 25 • a call to Reference, where the return value designates a value not equal to Initial_Value;
 - 26/2 • a call to Set_Value where the Val parameter is not equal to Initial_Value and the old attribute value is equal to Initial_Value;
 - 27 • a call to Set_Value where the Val parameter is not equal to Initial_Value and the old attribute value is not equal to Initial_Value.

Implementation Permissions

- 28 An implementation need not actually create the object corresponding to a task attribute until its value is set to something other than that of Initial_Value, or until Reference is called for the task attribute. Similarly, when the value of the attribute is to be reinitialized to that of Initial_Value, the object may instead be

finalized and its storage reclaimed, to be recreated when needed later. While the object does not exist, the function Value may simply return Initial_Value, rather than implicitly creating the object.

An implementation is allowed to place restrictions on the maximum number of attributes a task may have, the maximum size of each attribute, and the total storage size allocated for all the attributes of a task. 29

Implementation Advice

Some implementations are targeted to domains in which memory use at run time must be completely deterministic. For such implementations, it is recommended that the storage for task attributes will be pre-allocated statically and not from the heap. This can be accomplished by either placing restrictions on the number and the size of the attributes of a task, or by using the pre-allocated storage for the first N attribute objects, and the heap for the others. In the latter case, N should be documented. 30/2

Finalization of task attributes and reclamation of associated storage should be performed as soon as possible after task termination. 30.1/2

NOTES

12 An attribute always exists (after instantiation), and has the initial value. It need not occupy memory until the first operation that potentially changes the attribute value. The same holds true after Reinitialize. 31

13 The result of the Reference function should be used with care; it is always safe to use that result in the task body whose attribute is being accessed. However, when the result is being used by another task, the programmer must make sure that the task whose attribute is being accessed is not yet terminated. Failing to do so could make the program execution erroneous. 32

This paragraph was deleted.

33/2

C.7.3 The Package Task_Termination

Static Semantics

The following language-defined library package exists:

```
with Ada.Task_Identification;
with Ada.Exceptions;
package Ada.Task_Termination is
    pragma Preelaborate(Task_Termination);
    type Cause_Of_Termination is (Normal, Abnormal, Unhandled_Exception); 1/2
    type Termination_Handler is access protected procedure; 2/2
        (Cause : in Cause_Of_Termination;
         T     : in Ada.Task_Identification.Task_Id;
         X     : in Ada.Exceptions.Exception_Occurrence);
    procedure Set_Dependents_Fallback_Handler; 3/2
        (Handler: in Termination_Handler);
    function Current_Task_Fallback_Handler return Termination_Handler; 4/2
    procedure Set_Specific_Handler; 5/2
        (T      : in Ada.Task_Identification.Task_Id;
         Handler : in Termination_Handler);
    function Specific_Handler (T : Ada.Task_Identification.Task_Id) 6/2
        return Termination_Handler;
end Ada.Task_Termination; 7/2
```

Dynamic Semantics

The type Termination_Handler identifies a protected procedure to be executed by the implementation when a task terminates. Such a protected procedure is called a *handler*. In all cases T identifies the task that is terminating. If the task terminates due to completing the last statement of its body, or as a result of waiting on a terminate alternative, then Cause is set to Normal and X is set to Null_Occurrence. If the task

8/2

terminates because it is being aborted, then Cause is set to Abnormal and X is set to Null_Occurrence. If the task terminates because of an exception raised by the execution of its `task_body`, then Cause is set to Unhandled_Exception and X is set to the associated exception occurrence.

- 9/2 Each task has two termination handlers, a *fall-back handler* and a *specific handler*. The specific handler applies only to the task itself, while the fall-back handler applies only to the dependent tasks of the task. A handler is said to be *set* if it is associated with a non-null value of type `Termination_Handler`, and *cleared* otherwise. When a task is created, its specific handler and fall-back handler are cleared.
- 10/2 The procedure `Set_Dependents_Fallback_Handler` changes the fall-back handler for the calling task; if Handler is `null`, that fall-back handler is cleared, otherwise it is set to be Handler.`all`. If a fall-back handler had previously been set it is replaced.
- 11/2 The function `Current_Task_Fallback_Handler` returns the fall-back handler that is currently set for the calling task, if one is set; otherwise it returns `null`.
- 12/2 The procedure `Set_Specific_Handler` changes the specific handler for the task identified by T; if Handler is `null`, that specific handler is cleared, otherwise it is set to be Handler.`all`. If a specific handler had previously been set it is replaced.
- 13/2 The function `Specific_Handler` returns the specific handler that is currently set for the task identified by T, if one is set; otherwise it returns `null`.
- 14/2 As part of the finalization of a `task_body`, after performing the actions specified in 7.6 for finalization of a master, the specific handler for the task, if one is set, is executed. If the specific handler is cleared, a search for a fall-back handler proceeds by recursively following the master relationship for the task. If a task is found whose fall-back handler is set, that handler is executed; otherwise, no handler is executed.
- 15/2 For `Set_Specific_Handler` or `Specific_Handler`, `Tasking_Error` is raised if the task identified by T has already terminated. `Program_Error` is raised if the value of T is `Ada.Task_Identification.Null_Task_Id`.
- 16/2 An exception propagated from a handler that is invoked as part of the termination of a task has no effect.

Erroneous Execution

- 17/2 For a call of `Set_Specific_Handler` or `Specific_Handler`, if the task identified by T no longer exists, the execution of the program is erroneous.

Annex D (normative) Real-Time Systems

This Annex specifies additional characteristics of Ada implementations intended for real-time systems software. To conform to this Annex, an implementation shall also conform to the Systems Programming Annex.

Metrics

The metrics are documentation requirements; an implementation shall document the values of the language-defined metrics for at least one configuration of hardware or an underlying system supported by the implementation, and shall document the details of that configuration.

The metrics do not necessarily yield a simple number. For some, a range is more suitable, for others a formula dependent on some parameter is appropriate, and for others, it may be more suitable to break the metric into several cases. Unless specified otherwise, the metrics in this annex are expressed in processor clock cycles. For metrics that require documentation of an upper bound, if there is no upper bound, the implementation shall report that the metric is unbounded.

NOTES

1 The specification of the metrics makes a distinction between upper bounds and simple execution times. Where something is just specified as “the execution time of” a piece of code, this leaves one the freedom to choose a nonpathological case. This kind of metric is of the form “there exists a program such that the value of the metric is V”. Conversely, the meaning of upper bounds is “there is no program such that the value of the metric is greater than V”. This kind of metric can only be partially tested, by finding the value of V for one or more test programs.

2 The metrics do not cover the whole language; they are limited to features that are specified in Annex C, “Systems Programming” and in this Annex. The metrics are intended to provide guidance to potential users as to whether a particular implementation of such a feature is going to be adequate for a particular real-time application. As such, the metrics are aimed at known implementation choices that can result in significant performance differences.

3 The purpose of the metrics is not necessarily to provide fine-grained quantitative results or to serve as a comparison between different implementations on the same or different platforms. Instead, their goal is rather qualitative; to define a standard set of approximate values that can be measured and used to estimate the general suitability of an implementation, or to evaluate the comparative utility of certain features of an implementation for a particular real-time application.

D.1 Task Priorities

This clause specifies the priority model for real-time systems. In addition, the methods for specifying priorities are defined.

Syntax

The form of a **pragma Priority** is as follows:

pragma Priority(expression);

The form of a **pragma Interrupt_Priority** is as follows:

pragma Interrupt_Priority[(expression)];

Name Resolution Rules

The expected type for the expression in a Priority or Interrupt_Priority pragma is Integer.

Legality Rules

- 7 A Priority pragma is allowed only immediately within a `task_definition`, a `protected_definition`, or the `declarative_part` of a `subprogram_body`. An `Interrupt_Priority` pragma is allowed only immediately within a `task_definition` or a `protected_definition`. At most one such pragma shall appear within a given construct.
- 8 For a Priority pragma that appears in the `declarative_part` of a `subprogram_body`, the `expression` shall be static, and its value shall be in the range of `System.Priority`.

Static Semantics

- 9 The following declarations exist in package `System`:

```
10    subtype Any_Priority is Integer range implementation-defined;
     subtype Priority is Any_Priority
          range Any_Priority'First .. implementation-defined;
     subtype Interrupt_Priority is Any_Priority
          range Priority'Last+1 .. Any_Priority'Last;
11    Default_Priority : constant Priority := (Priority'First + Priority'Last)/2;
```

- 12 The full range of priority values supported by an implementation is specified by the subtype `Any_Priority`. The subrange of priority values that are high enough to require the blocking of one or more interrupts is specified by the subtype `Interrupt_Priority`. The subrange of priority values below `System.Interrupt_Priority'First` is specified by the subtype `System.Priority`.

- 13 The priority specified by a Priority or `Interrupt_Priority` pragma is the value of the `expression` in the pragma, if any. If there is no `expression` in an `Interrupt_Priority` pragma, the priority value is `Interrupt_Priority'Last`.

Dynamic Semantics

- 14 A Priority pragma has no effect if it occurs in the `declarative_part` of the `subprogram_body` of a subprogram other than the main subprogram.

- 15 A *task priority* is an integer value that indicates a degree of urgency and is the basis for resolving competing demands of tasks for resources. Unless otherwise specified, whenever tasks compete for processors or other implementation-defined resources, the resources are allocated to the task with the highest priority value. The *base priority* of a task is the priority with which it was created, or to which it was later set by `Dynamic_Priorities.Set_Priority` (see D.5). At all times, a task also has an *active priority*, which generally reflects its base priority as well as any priority it inherits from other sources. *Priority inheritance* is the process by which the priority of a task or other entity (e.g. a protected object; see D.3) is used in the evaluation of another task's active priority.

- 16 The effect of specifying such a pragma in a `protected_definition` is discussed in D.3.
- 17 The `expression` in a Priority or `Interrupt_Priority` pragma that appears in a `task_definition` is evaluated for each task object (see 9.1). For a Priority pragma, the value of the `expression` is converted to the subtype `Priority`; for an `Interrupt_Priority` pragma, this value is converted to the subtype `Any_Priority`. The priority value is then associated with the task object whose `task_definition` contains the pragma.
- 18 Likewise, the priority value is associated with the environment task if the pragma appears in the `declarative_part` of the main subprogram.
- 19 The initial value of a task's base priority is specified by default or by means of a Priority or `Interrupt_Priority` pragma. After a task is created, its base priority can be changed only by a call to `Dynamic_Priorities.Set_Priority` (see D.5). The initial base priority of a task in the absence of a pragma is

the base priority of the task that creates it at the time of creation (see 9.1). If a pragma Priority does not apply to the main subprogram, the initial base priority of the environment task is System.Default_Priority. The task's active priority is used when the task competes for processors. Similarly, the task's active priority is used to determine the task's position in any queue when Priority_Queuing is specified (see D.4).

At any time, the active priority of a task is the maximum of all the priorities the task is inheriting at that instant. For a task that is not held (see D.11), its base priority is a source of priority inheritance unless otherwise specified for a particular task dispatching policy. Other sources of priority inheritance are specified under the following conditions:

- During activation, a task being activated inherits the active priority that its activator (see 9.2) had at the time the activation was initiated. 21/1
- During rendezvous, the task accepting the entry call inherits the priority of the entry call (see 9.5.3 and D.4). 22/1
- During a protected action on a protected object, a task inherits the ceiling priority of the protected object (see 9.5 and D.3). 23

In all of these cases, the priority ceases to be inherited as soon as the condition calling for the inheritance no longer exists. 24

Implementation Requirements

The range of System.Interrupt_Priority shall include at least one value. 25

The range of System.Priority shall include at least 30 values. 26

NOTES

4 The priority expression can include references to discriminants of the enclosing type. 27

5 It is a consequence of the active priority rules that at the point when a task stops inheriting a priority from another source, its active priority is re-evaluated. This is in addition to other instances described in this Annex for such re-evaluation. 28

6 An implementation may provide a non-standard mode in which tasks inherit priorities under conditions other than those specified above. 29

D.2 Priority Scheduling

This clause describes the rules that determine which task is selected for execution when more than one task is ready (see 9). 1/2

D.2.1 The Task Dispatching Model

The task dispatching model specifies task scheduling, based on conceptual priority-ordered ready queues. 1/2

Static Semantics

The following language-defined library package exists: 1.1/2

```
package Ada.Dispatching is
  pragma Pure(Dispatching);
  Dispatching_Policy_Error : exception;
end Ada.Dispatching;
```

Dispatching serves as the parent of other language-defined library units concerned with task dispatching. 1.3/2

Dynamic Semantics

- 2/2 A task can become a *running task* only if it is ready (see 9) and the execution resources required by that task are available. Processors are allocated to tasks based on each task's active priority.
- 3 It is implementation defined whether, on a multiprocessor, a task that is waiting for access to a protected object keeps its processor busy.
- 4/2 *Task dispatching* is the process by which one ready task is selected for execution on a processor. This selection is done at certain points during the execution of a task called *task dispatching points*. A task reaches a task dispatching point whenever it becomes blocked, and when it terminates. Other task dispatching points are defined throughout this Annex for specific policies.
- 5/2 *Task dispatching policies* are specified in terms of conceptual *ready queues* and task states. A ready queue is an ordered list of ready tasks. The first position in a queue is called the *head of the queue*, and the last position is called the *tail of the queue*. A task is *ready* if it is in a ready queue, or if it is running. Each processor has one ready queue for each priority value. At any instant, each ready queue of a processor contains exactly the set of tasks of that priority that are ready for execution on that processor, but are not running on any processor; that is, those tasks that are ready, are not running on any processor, and can be executed using that processor and other available resources. A task can be on the ready queues of more than one processor.
- 6/2 Each processor also has one *running task*, which is the task currently being executed by that processor. Whenever a task running on a processor reaches a task dispatching point it goes back to one or more ready queues; a task (possibly the same task) is then selected to run on that processor. The task selected is the one at the head of the highest priority nonempty ready queue; this task is then removed from all ready queues to which it belongs.
- 7/2 *This paragraph was deleted.*
- 8/2 *This paragraph was deleted.*

Implementation Permissions

- 9/2 An implementation is allowed to define additional resources as execution resources, and to define the corresponding allocation policies for them. Such resources may have an implementation-defined effect on task dispatching.
- 10 An implementation may place implementation-defined restrictions on tasks whose active priority is in the Interrupt_Priority range.
- 10.1/2 For optimization purposes, an implementation may alter the points at which task dispatching occurs, in an implementation-defined manner. However, a `delay_statement` always corresponds to at least one task dispatching point.

NOTES

- 11 7 Section 9 specifies under which circumstances a task becomes ready. The ready state is affected by the rules for task activation and termination, delay statements, and entry calls. When a task is not ready, it is said to be blocked.
- 12 8 An example of a possible implementation-defined execution resource is a page of physical memory, which needs to be loaded with a particular page of virtual memory before a task can continue execution.
- 13 9 The ready queues are purely conceptual; there is no requirement that such lists physically exist in an implementation.
- 14 10 While a task is running, it is not on any ready queue. Any time the task that is running on a processor is added to a ready queue, a new running task is selected for that processor.
- 15 11 In a multiprocessor system, a task can be on the ready queues of more than one processor. At the extreme, if several processors share the same set of ready tasks, the contents of their ready queues is identical, and so they can be viewed as

sharing one ready queue, and can be implemented that way. Thus, the dispatching model covers multiprocessors where dispatching is implemented using a single ready queue, as well as those with separate dispatching domains.

12 The priority of a task is determined by rules specified in this subclause, and under D.1, “Task Priorities”, D.3, “Priority Ceiling Locking”, and D.5, “Dynamic Priorities”. 16

13 The setting of a task's base priority as a result of a call to Set_Priority does not always take effect immediately when Set_Priority is called. The effect of setting the task's base priority is deferred while the affected task performs a protected action. 17/2

D.2.2 Task Dispatching Pragmas

This clause allows a single task dispatching policy to be defined for all priorities, or the range of priorities to be split into subranges that are assigned individual dispatching policies. 0.1/2

Syntax

The form of a **pragma** Task_Dispatching_Policy is as follows: 1

pragma Task_Dispatching_Policy(*policy_identifier*); 2

The form of a **pragma** Priority_Specific_Dispatching is as follows: 2.1/2

pragma Priority_Specific_Dispatching (2.2/2
policy_identifier,*first_priority_expression*,*last_priority_expression*);

Name Resolution Rules

The expected type for *first_priority_expression* and *last_priority_expression* is Integer. 2.3/2

Legality Rules

The *policy_identifier* used in a **pragma** Task_Dispatching_Policy shall be the name of a task dispatching policy. 3/2

The *policy_identifier* used in a **pragma** Priority_Specific_Dispatching shall be the name of a task dispatching policy. 3.1/2

Both *first_priority_expression* and *last_priority_expression* shall be static expressions in the range of System.Any_Priority; *last_priority_expression* shall have a value greater than or equal to *first_priority_expression*. 3.2/2

Static Semantics

Pragma Task_Dispatching_Policy specifies the single task dispatching policy. 3.3/2

Pragma Priority_Specific_Dispatching specifies the task dispatching policy for the specified range of priorities. Tasks with base priorities within the range of priorities specified in a Priority_Specific_Dispatching pragma have their active priorities determined according to the specified dispatching policy. Tasks with active priorities within the range of priorities specified in a Priority_Specific_Dispatching pragma are dispatched according to the specified dispatching policy. 3.4/2

If a partition contains one or more Priority_Specific_Dispatching pragmas the dispatching policy for priorities not covered by any Priority_Specific_Dispatching pragmas is FIFO_Within_Priorities. 3.5/2

Post-Compilation Rules

A Task_Dispatching_Policy pragma is a configuration pragma. A Priority_Specific_Dispatching pragma is a configuration pragma. 4/2

- 4.1/2 The priority ranges specified in more than one Priority_Specific_Dispatching pragma within the same partition shall not be overlapping.
- 4.2/2 If a partition contains one or more Priority_Specific_Dispatching pragmas it shall not contain a Task_Dispatching_Policy pragma.
- 5/2 *This paragraph was deleted.*

Dynamic Semantics

- 6/2 A *task dispatching policy* specifies the details of task dispatching that are not covered by the basic task dispatching model. These rules govern when tasks are inserted into and deleted from the ready queues. A single task dispatching policy is specified by a Task_Dispatching_Policy pragma. Pragma Priority_Specific_Dispatching assigns distinct dispatching policies to subranges of System.Any_Priority.
- 6.1/2 If neither *pragma* applies to any of the program units comprising a partition, the task dispatching policy for that partition is unspecified.
- 6.2/2 If a partition contains one or more Priority_Specific_Dispatching pragmas a task dispatching point occurs for the currently running task of a processor whenever there is a non-empty ready queue for that processor with a higher priority than the priority of the running task.
- 6.3/2 A task that has its base priority changed may move from one dispatching policy to another. It is immediately subject to the new dispatching policy.

Paragraphs 7 through 13 were moved to D.2.3.

Implementation Requirements

- 13.1/2 An implementation shall allow, for a single partition, both the locking policy (see D.3) to be specified as Ceiling_Locking and also one or more Priority_Specific_Dispatching pragmas to be given.

Documentation Requirements

Paragraphs 14 through 16 were moved to D.2.3.

Implementation Permissions

- 17/2 Implementations are allowed to define other task dispatching policies, but need not support more than one task dispatching policy per partition.
- 18/2 An implementation need not support *pragma* Priority_Specific_Dispatching if it is infeasible to support it in the target environment.

NOTES

Paragraphs 19 through 21 were deleted.

D.2.3 Preemptive Dispatching

- 1/2 This clause defines a preemptive task dispatching policy.

Static Semantics

- 2/2 The *policy_identifier* FIFO_Within_Priorities is a task dispatching policy.

Dynamic Semantics

- 3/2 When FIFO_Within_Priorities is in effect, modifications to the ready queues occur only as follows:

- When a blocked task becomes ready, it is added at the tail of the ready queue for its active priority. 4/2
- When the active priority of a ready task that is not running changes, or the setting of its base priority takes effect, the task is removed from the ready queue for its old active priority and is added at the tail of the ready queue for its new active priority, except in the case where the active priority is lowered due to the loss of inherited priority, in which case the task is added at the head of the ready queue for its new active priority. 5/2
- When the setting of the base priority of a running task takes effect, the task is added to the tail of the ready queue for its active priority. 6/2
- When a task executes a `delay_statement` that does not result in blocking, it is added to the tail of the ready queue for its active priority. 7/2

Each of the events specified above is a task dispatching point (see D.2.1). 8/2

A task dispatching point occurs for the currently running task of a processor whenever there is a nonempty ready queue for that processor with a higher priority than the priority of the running task. The currently running task is said to be *preempted* and it is added at the head of the ready queue for its active priority. 9/2

Implementation Requirements

An implementation shall allow, for a single partition, both the task dispatching policy to be specified as `FIFO_Within_Priorities` and also the locking policy (see D.3) to be specified as `Ceiling_Locking`. 10/2

Documentation Requirements

Priority inversion is the duration for which a task remains at the head of the highest priority nonempty ready queue while the processor executes a lower priority task. The implementation shall document: 11/2

- The maximum priority inversion a user task can experience due to activity of the implementation (on behalf of lower priority tasks), and 12/2
- whether execution of a task can be preempted by the implementation processing of delay expirations for lower priority tasks, and if so, for how long. 13/2

NOTES

14 If the active priority of a running task is lowered due to loss of inherited priority (as it is on completion of a protected operation) and there is a ready task of the same active priority that is not running, the running task continues to run (provided that there is no higher priority task). 14/2

15 Setting the base priority of a ready task causes the task to move to the tail of the queue for its active priority, regardless of whether the active priority of the task actually changes. 15/2

D.2.4 Non-Preemptive Dispatching

This clause defines a non-preemptive task dispatching policy. 1/2

Static Semantics

The `policy_identifier Non_Preemptive_FIFO_Within_Priorities` is a task dispatching policy. 2/2

Legality Rules

`Non_Preemptive_FIFO_Within_Priorities` shall not be specified as the `policy_identifier` of pragma `Priority_Specific_Dispatching` (see D.2.2). 3/2

Dynamic Semantics

- 4/2 When Non_Preemptive_FIFO_Within_Priorities is in effect, modifications to the ready queues occur only as follows:
- 5/2 • When a blocked task becomes ready, it is added at the tail of the ready queue for its active priority.
 - 6/2 • When the active priority of a ready task that is not running changes, or the setting of its base priority takes effect, the task is removed from the ready queue for its old active priority and is added at the tail of the ready queue for its new active priority.
 - 7/2 • When the setting of the base priority of a running task takes effect, the task is added to the tail of the ready queue for its active priority.
 - 8/2 • When a task executes a `delay_statement` that does not result in blocking, it is added to the tail of the ready queue for its active priority.
- 9/2 For this policy, a non-blocking `delay_statement` is the only non-blocking event that is a task dispatching point (see D.2.1).

Implementation Requirements

- 10/2 An implementation shall allow, for a single partition, both the task dispatching policy to be specified as Non_Preemptive_FIFO_Within_Priorities and also the locking policy (see D.3) to be specified as Ceiling_Locking.

Implementation Permissions

- 11/2 Since implementations are allowed to round all ceiling priorities in subrange System.Priority to System.Priority'Last (see D.3), an implementation may allow a task to execute within a protected object without raising its active priority provided the associated protected unit does not contain pragma Interrupt_Priority, Interrupt_Handler, or Attach_Handler.

D.2.5 Round Robin Dispatching

- 1/2 This clause defines the task dispatching policy Round_Robin_Within_Priorities and the package Round_Robin.

Static Semantics

- 2/2 The `policy_identifier` Round_Robin_Within_Priorities is a task dispatching policy.

- 3/2 The following language-defined library package exists:

```
4/2   with System;
      with Ada.Real_Time;
      package Ada.Dispatching.Round_Robin is
        Default_Quantum : constant Ada.Real_Time.Time_Span :=
          implementation-defined;
        procedure Set_Quantum (Pri      : in System.Priority;
                               Quantum : in Ada.Real_Time.Time_Span);
        procedure Set_Quantum (Low, High : in System.Priority;
                               Quantum : in Ada.Real_Time.Time_Span);
        function Actual_Quantum (Pri : System.Priority) return
          Ada.Real_Time.Time_Span;
        function Is_Round_Robin (Pri : System.Priority) return Boolean;
      end Ada.Dispatching.Round_Robin;
```

When task dispatching policy Round_Robin_Within_Priorities is the single policy in effect for a partition, each task with priority in the range of System.Interrupt_Priority is dispatched according to policy FIFO_Within_Priorities.

5/2

Dynamic Semantics

The procedures Set_Quantum set the required Quantum value for a single priority level Pri or a range of priority levels Low .. High. If no quantum is set for a Round Robin priority level, Default_Quantum is used.

6/2

The function Actual_Quantum returns the actual quantum used by the implementation for the priority level Pri.

7/2

The function Is_Round_Robin returns True if priority Pri is covered by task dispatching policy Round_Robin_Within_Priorities; otherwise it returns False.

8/2

A call of Actual_Quantum or Set_Quantum raises exception Dispatching.Dispatching_Policy_Error if a predefined policy other than Round_Robin_Within_Priorities applies to the specified priority or any of the priorities in the specified range.

9/2

For Round_Robin_Within_Priorities, the dispatching rules for FIFO_Within_Priorities apply with the following additional rules:

10/2

- When a task is added or moved to the tail of the ready queue for its base priority, it has an execution time budget equal to the quantum for that priority level. This will also occur when a blocked task becomes executable again.
- When a task is preempted (by a higher priority task) and is added to the head of the ready queue for its priority level, it retains its remaining budget.
- While a task is executing, its budget is decreased by the amount of execution time it uses. The accuracy of this accounting is the same as that for execution time clocks (see D.14).
- When a task has exhausted its budget and is without an inherited priority (and is not executing within a protected operation), it is moved to the tail of the ready queue for its priority level. This is a task dispatching point.

11/2

12/2

13/2

14/2

Implementation Requirements

An implementation shall allow, for a single partition, both the task dispatching policy to be specified as Round_Robin_Within_Priorities and also the locking policy (see D.3) to be specified as Ceiling_Locking.

15/2

Documentation Requirements

An implementation shall document the quantum values supported.

16/2

An implementation shall document the accuracy with which it detects the exhaustion of the budget of a task.

17/2

NOTES

16 Due to implementation constraints, the quantum value returned by Actual_Quantum might not be identical to that set with Set_Quantum.

18/2

17 A task that executes continuously with an inherited priority will not be subject to round robin dispatching.

19/2

D.2.6 Earliest Deadline First Dispatching

- 1/2 The deadline of a task is an indication of the urgency of the task; it represents a point on an ideal physical time line. The deadline might affect how resources are allocated to the task.
- 2/2 This clause defines a package for representing the deadline of a task and a dispatching policy that defines Earliest Deadline First (EDF) dispatching. A pragma is defined to assign an initial deadline to a task.

Syntax

- 3/2 The form of a **pragma** `Relative_Deadline` is as follows:
- 4/2 **pragma** `Relative_Deadline (relative_deadline_expression);`

Name Resolution Rules

- 5/2 The expected type for `relative_deadline_expression` is `Real_Time.Time_Span`.

Legality Rules

- 6/2 A `Relative_Deadline` pragma is allowed only immediately within a `task_definition` or the `declarative_part` of a `subprogram_body`. At most one such pragma shall appear within a given construct.

Static Semantics

- 7/2 The `policy_identifier` `EDF_Across_Priorities` is a task dispatching policy.

- 8/2 The following language-defined library package exists:

```
9/2 with Ada.Real_Time;
with Ada.Task_Identification;
package Ada.Dispatching.EDF is
    subtype Deadline is Ada.Real_Time.Time;
    Default_Deadline : constant Deadline :=
        Ada.Real_Time.Time_Last;
    procedure Set_Deadline (D : in Deadline;
                           T : in Ada.Task_Identification.Task_Id := Ada.Task_Identification.Current_Task);
    procedure Delay_Until_And_Set_Deadline (
        Delay_Until_Time : in Ada.Real_Time.Time;
        Deadline_Offset : in Ada.Real_Time.Time_Span);
    function Get_Deadline (T : Ada.Task_Identification.Task_Id := Ada.Task_Identification.Current_Task) return Deadline;
end Ada.Dispatching.EDF;
```

Post-Compilation Rules

- 10/2 If the `EDF_Across_Priorities` policy is specified for a partition, then the `Ceiling_Locking` policy (see D.3) shall also be specified for the partition.
- 11/2 If the `EDF_Across_Priorities` policy appears in a `Priority_Specific_Dispatching` pragma (see D.2.2) in a partition, then the `Ceiling_Locking` policy (see D.3) shall also be specified for the partition.

Dynamic Semantics

- 12/2 A `Relative_Deadline` pragma has no effect if it occurs in the `declarative_part` of the `subprogram_body` of a subprogram other than the main subprogram.
- 13/2 The initial absolute deadline of a task containing `pragma Relative_Deadline` is the value of `Real_Time.Clock + relative_deadline_expression`, where the call of `Real_Time.Clock` is made between task creation and the start of its activation. If there is no `Relative_Deadline` pragma then the initial absolute

deadline of a task is the value of `Default_Deadline`. The environment task is also given an initial deadline by this rule.

The procedure `Set_Deadline` changes the absolute deadline of the task to `D`. The function `Get_Deadline` returns the absolute deadline of the task. 14/2

The procedure `Delay_Until_And_Set_Deadline` delays the calling task until time `Delay_Until_Time`. When the task becomes runnable again it will have deadline `Delay_Until_Time + Deadline_Offset`. 15/2

On a system with a single processor, the setting of the deadline of a task to the new value occurs immediately at the first point that is outside the execution of a protected action. If the task is currently on a ready queue it is removed and re-entered on to the ready queue determined by the rules defined below. 16/2

When `EDF_Across_Priorities` is specified for priority range `Low..High` all ready queues in this range are ordered by deadline. The task at the head of a queue is the one with the earliest deadline. 17/2

A task dispatching point occurs for the currently running task `T` to which policy `EDF_Across_Priorities` applies: 18/2

- when a change to the deadline of `T` occurs; 19/2
- there is a task on the ready queue for the active priority of `T` with a deadline earlier than the deadline of `T`; or 20/2
- there is a non-empty ready queue for that processor with a higher priority than the active priority of the running task. 21/2

In these cases, the currently running task is said to be preempted and is returned to the ready queue for its active priority. 22/2

For a task `T` to which policy `EDF_Across_Priorities` applies, the base priority is not a source of priority inheritance; the active priority when first activated or while it is blocked is defined as the maximum of the following: 23/2

- the lowest priority in the range specified as `EDF_Across_Priorities` that includes the base priority of `T`; 24/2
- the priorities, if any, currently inherited by `T`; 25/2
- the highest priority `P`, if any, less than the base priority of `T` such that one or more tasks are executing within a protected object with ceiling priority `P` and task `T` has an earlier deadline than all such tasks. 26/2

When a task `T` is first activated or becomes unblocked, it is added to the ready queue corresponding to this active priority. Until it becomes blocked again, the active priority of `T` remains no less than this value; it will exceed this value only while it is inheriting a higher priority. 27/2

When the setting of the base priority of a ready task takes effect and the new priority is in a range specified as `EDF_Across_Priorities`, the task is added to the ready queue corresponding to its new active priority, as determined above. 28/2

For all the operations defined in `Dispatching.EDF`, `Tasking_Error` is raised if the task identified by `T` has terminated. `Program_Error` is raised if the value of `T` is `Null_Task_Id`. 29/2

Bounded (Run-Time) Errors

If `EDF_Across_Priorities` is specified for priority range `Low..High`, it is a bounded error to declare a protected object with ceiling priority `Low` or to assign the value `Low` to attribute `'Priority`. In either case either `Program_Error` is raised or the ceiling of the protected object is assigned the value `Low+1`. 30/2

Erroneous Execution

- 31/2 If a value of Task_Id is passed as a parameter to any of the subprograms of this package and the corresponding task object no longer exists, the execution of the program is erroneous.

Documentation Requirements

- 32/2 On a multiprocessor, the implementation shall document any conditions that cause the completion of the setting of the deadline of a task to be delayed later than what is specified for a single processor.

NOTES

- 33/2 18 If two adjacent priority ranges, $A..B$ and $B+1..C$ are specified to have policy EDF_Across_Priorities then this is not equivalent to this policy being specified for the single range, $A..C$.

- 34/2 19 The above rules implement the preemption-level protocol (also called Stack Resource Policy protocol) for resource sharing under EDF dispatching. The preemption-level for a task is denoted by its base priority. The definition of a ceiling preemption-level for a protected object follows the existing rules for ceiling locking.

D.3 Priority Ceiling Locking

- 1 This clause specifies the interactions between priority task scheduling and protected object ceilings. This interaction is based on the concept of the *ceiling priority* of a protected object.

Syntax

2 The form of a **pragma** Locking_Policy is as follows:

3 **pragma** Locking_Policy(*policy_identifier*);

Legality Rules

- 4 The *policy_identifier* shall either be Ceiling_Locking or an implementation-defined identifier.

Post-Compilation Rules

- 5 A Locking_Policy pragma is a configuration pragma.

Dynamic Semantics

- 6/2 A locking policy specifies the details of protected object locking. All protected objects have a priority. The locking policy specifies the meaning of the priority of a protected object, and the relationships between these priorities and task priorities. In addition, the policy specifies the state of a task when it executes a protected action, and how its active priority is affected by the locking. The *locking policy* is specified by a Locking_Policy pragma. For implementation-defined locking policies, the meaning of the priority of a protected object is implementation defined. If no Locking_Policy pragma applies to any of the program units comprising a partition, the locking policy for that partition, as well as the meaning of the priority of a protected object, are implementation defined.

- 6.1/2 The expression of a Priority or Interrupt_Priority pragma (see D.1) is evaluated as part of the creation of the corresponding protected object and converted to the subtype System.Any_Priority or System.Interrupt_Priority, respectively. The value of the expression is the initial priority of the corresponding protected object. If no Priority or Interrupt_Priority pragma applies to a protected object, the initial priority is specified by the locking policy.

- 7 There is one predefined locking policy, Ceiling_Locking; this policy is defined as follows:

- 8/2 • Every protected object has a *ceiling priority*, which is determined by either a Priority or Interrupt_Priority pragma as defined in D.1, or by assignment to the Priority attribute as described in D.5.2. The ceiling priority of a protected object (or ceiling, for short) is an upper

bound on the active priority a task can have when it calls protected operations of that protected object.

- The initial ceiling priority of a protected object is equal to the initial priority for that object. 9/2
- If an Interrupt_Handler or Attach_Handler pragma (see C.3.1) appears in a protected_definition without an Interrupt_Priority pragma, the initial priority of protected objects of that type is implementation defined, but in the range of the subtype System.Interrupt_Priority. 10/2
- If no pragma Priority, Interrupt_Priority, Interrupt_Handler, or Attach_Handler is specified in the protected_definition, then the initial priority of the corresponding protected object is System.Priority'Last. 11/2
- While a task executes a protected action, it inherits the ceiling priority of the corresponding protected object. 12
- When a task calls a protected operation, a check is made that its active priority is not higher than the ceiling of the corresponding protected object; Program_Error is raised if this check fails. 13

Bounded (Run-Time) Errors

Following any change of priority, it is a bounded error for the active priority of any task with a call queued on an entry of a protected object to be higher than the ceiling priority of the protected object. In this case one of the following applies:

- at any time prior to executing the entry body Program_Error is raised in the calling task; 13.2/2
- when the entry is open the entry body is executed at the ceiling priority of the protected object; 13.3/2
- when the entry is open the entry body is executed at the ceiling priority of the protected object and then Program_Error is raised in the calling task; or 13.4/2
- when the entry is open the entry body is executed at the ceiling priority of the protected object that was in effect when the entry call was queued. 13.5/2

Implementation Permissions

The implementation is allowed to round all ceilings in a certain subrange of System.Priority or System.Interrupt_Priority up to the top of that subrange, uniformly. 14

Implementations are allowed to define other locking policies, but need not support more than one locking policy per partition. 15/2

Since implementations are allowed to place restrictions on code that runs at an interrupt-level active priority (see C.3.1 and D.2.1), the implementation may implement a language feature in terms of a protected object with an implementation-defined ceiling, but the ceiling shall be no less than Priority'Last. 16

Implementation Advice

The implementation should use names that end with “_Locking” for implementation-defined locking policies. 17

NOTES

20 While a task executes in a protected action, it can be preempted only by tasks whose active priorities are higher than the ceiling priority of the protected object. 18

21 If a protected object has a ceiling priority in the range of Interrupt_Priority, certain interrupts are blocked while protected actions of that object execute. In the extreme, if the ceiling is Interrupt_Priority'Last, all blockable interrupts are blocked during that time. 19

22 The ceiling priority of a protected object has to be in the Interrupt_Priority range if one of its procedures is to be used as an interrupt handler (see C.3). 20

- 21 23 When specifying the ceiling of a protected object, one should choose a value that is at least as high as the highest active priority at which tasks can be executing when they call protected operations of that object. In determining this value the following factors, which can affect active priority, should be considered: the effect of Set_Priority, nested protected operations, entry calls, task activation, and other implementation-defined factors.
- 22 24 Attaching a protected procedure whose ceiling is below the interrupt hardware priority to an interrupt causes the execution of the program to be erroneous (see C.3.1).
- 23 25 On a single processor implementation, the ceiling priority rules guarantee that there is no possibility of deadlock involving only protected subprograms (excluding the case where a protected operation calls another protected operation on the same protected object).

D.4 Entry Queuing Policies

- 1/1 This clause specifies a mechanism for a user to choose an entry *queuing policy*. It also defines two such policies. Other policies are implementation defined.

Syntax

- 2 The form of a **pragma** Queueing_Policy is as follows:
- 3 **pragma** Queueing_Policy(*policy_identifier*);

Legality Rules

- 4 The *policy_identifier* shall be either FIFO_Queueing, Priority_Queueing or an implementation-defined identifier.

Post-Compilation Rules

- 5 A Queueing_Policy pragma is a configuration pragma.

Dynamic Semantics

- 6 A *queuing policy* governs the order in which tasks are queued for entry service, and the order in which different entry queues are considered for service. The queuing policy is specified by a Queueing_Policy pragma.
- 7/2 Two queuing policies, FIFO_Queueing and Priority_Queueing, are language defined. If no Queueing_Policy pragma applies to any of the program units comprising the partition, the queuing policy for that partition is FIFO_Queueing. The rules for this policy are specified in 9.5.3 and 9.7.1.
- 8 The Priority_Queueing policy is defined as follows:
- 9 • The calls to an entry (including a member of an entry family) are queued in an order consistent with the priorities of the calls. The *priority of an entry call* is initialized from the active priority of the calling task at the time the call is made, but can change later. Within the same priority, the order is consistent with the calling (or requeueing, or priority setting) time (that is, a FIFO order).
 - 10/1 • After a call is first queued, changes to the active priority of a task do not affect the priority of the call, unless the base priority of the task is set while the task is blocked on an entry call.
 - 11 • When the base priority of a task is set (see D.5), if the task is blocked on an entry call, and the call is queued, the priority of the call is updated to the new active priority of the calling task. This causes the call to be removed from and then reinserted in the queue at the new active priority.
 - 12 • When more than one condition of an entry_barrier of a protected object becomes True, and more than one of the respective queues is nonempty, the call with the highest priority is selected. If more than one such call has the same priority, the call that is queued on the entry whose

declaration is first in textual order in the `protected_definition` is selected. For members of the same entry family, the one with the lower family index is selected.

- If the expiration time of two or more open `delay_alternatives` is the same and no other `accept_alternatives` are open, the `sequence_of_statements` of the `delay_alternative` that is first in textual order in the `selective_accept` is executed. 13
- When more than one alternative of a `selective_accept` is open and has queued calls, an alternative whose queue has the highest-priority call at its head is selected. If two or more open alternatives have equal-priority queued calls, then a call on the entry in the `accept_alternative` that is first in textual order in the `selective_accept` is selected. 14

Implementation Permissions

Implementations are allowed to define other queuing policies, but need not support more than one queuing policy per partition. 15/2

Implementations are allowed to defer the reordering of entry queues following a change of base priority of a task blocked on the entry call if it is not practical to reorder the queue immediately. 15.1/2

Implementation Advice

The implementation should use names that end with “_Queuing” for implementation-defined queuing policies. 16

D.5 Dynamic Priorities

This clause describes how the priority of an entity can be modified or queried at run time. 1/2

D.5.1 Dynamic Priorities for Tasks

This clause describes how the base priority of a task can be modified or queried at run time. 1

Static Semantics

The following language-defined library package exists: 2

```
with System;
with Ada.Task_Identification; -- See C.7.1
package Ada.Dynamic_Priorities is
  pragma Preelaborate(Dynamic_Priorities);
  procedure Set_Priority(Priority : in System.Any_Priority;
                        T : in Ada.Task_Identification.Task_Id := Ada.Task_Identification.Current_Task);
  function Get_Priority (T : Ada.Task_Identification.Task_Id := Ada.Task_Identification.Current_Task)
                        return System.Any_Priority;
end Ada.Dynamic_Priorities; 6
```

Dynamic Semantics

The procedure `Set_Priority` sets the base priority of the specified task to the specified Priority value. `Set_Priority` has no effect if the task is terminated. 7

The function `Get_Priority` returns T's current base priority. `Tasking_Error` is raised if the task is terminated. 8

`Program_Error` is raised by `Set_Priority` and `Get_Priority` if T is equal to `Null_Task_Id`. 9

- 10/2 On a system with a single processor, the setting of the base priority of a task T to the new value occurs immediately at the first point when T is outside the execution of a protected action.

Bounded (Run-Time) Errors

- 11/2 *This paragraph was deleted.*

Erroneous Execution

- 12 If any subprogram in this package is called with a parameter T that specifies a task object that no longer exists, the execution of the program is erroneous.

Documentation Requirements

- 12.1/2 On a multiprocessor, the implementation shall document any conditions that cause the completion of the setting of the priority of a task to be delayed later than what is specified for a single processor.

Metrics

- 13 The implementation shall document the following metric:

- 14 • The execution time of a call to `Set_Priority`, for the nonpreempting case, in processor clock cycles. This is measured for a call that modifies the priority of a ready task that is not running (which cannot be the calling one), where the new base priority of the affected task is lower than the active priority of the calling task, and the affected task is not on any entry queue and is not executing a protected operation.

NOTES

- 15/2 26 Setting a task's base priority affects task dispatching. First, it can change the task's active priority. Second, under the `FIFO_Within_Priorities` policy it always causes the task to move to the tail of the ready queue corresponding to its active priority, even if the new base priority is unchanged.

- 16 27 Under the priority queuing policy, setting a task's base priority has an effect on a queued entry call if the task is blocked waiting for the call. That is, setting the base priority of a task causes the priority of a queued entry call from that task to be updated and the call to be removed and then reinserted in the entry queue at the new priority (see D.4), unless the call originated from the `triggering_statement` of an `asynchronous_select`.

- 17 28 The effect of two or more `Set_Priority` calls executed in parallel on the same task is defined as executing these calls in some serial order.

- 18 29 The rule for when `Tasking_Error` is raised for `Set_Priority` or `Get_Priority` is different from the rule for when `Tasking_Error` is raised on an entry call (see 9.5.3). In particular, setting or querying the priority of a completed or an abnormal task is allowed, so long as the task is not yet terminated.

- 19 30 Changing the priorities of a set of tasks can be performed by a series of calls to `Set_Priority` for each task separately. For this to work reliably, it should be done within a protected operation that has high enough ceiling priority to guarantee that the operation completes without being preempted by any of the affected tasks.

D.5.2 Dynamic Priorities for Protected Objects

This clause specifies how the priority of a protected object can be modified or queried at run time.

1/2

Static Semantics

The following attribute is defined for a prefix P that denotes a protected object:

2/2

P'Priority Denotes a non-aliased component of the protected object P. This component is of type System.Any_Priority and its value is the priority of P. P'Priority denotes a variable if and only if P denotes a variable. A reference to this attribute shall appear only within the body of P.

3/2

The initial value of this attribute is the initial value of the priority of the protected object, and can be changed by an assignment.

4/2

Dynamic Semantics

If the locking policy Ceiling_Locking (see D.3) is in effect then the ceiling priority of a protected object P is set to the value of P'Priority at the end of each protected action of P.

5/2

If the locking policy Ceiling_Locking is in effect, then for a protected object P with either an Attach_Handler or Interrupt_Handler pragma applying to one of its procedures, a check is made that the value to be assigned to P'Priority is in the range System.Interrupt_Priority. If the check fails, Program_Error is raised.

6/2

Metrics

The implementation shall document the following metric:

7/2

- The difference in execution time of calls to the following procedures in protected object P:

8/2

```

protected P is
    procedure Do_Not_Set_Ceiling (Pr : System.Any_Priority);
    procedure Set_Ceiling (Pr : System.Any_Priority);
end P;

protected body P is
    procedure Do_Not_Set_Ceiling (Pr : System.Any_Priority) is
    begin
        null;
    end;
    procedure Set_Ceiling (Pr : System.Any_Priority) is
    begin
        P'Priority := Pr;
    end;
end P;

```

9/2

10/2

NOTES

31 Since P'Priority is a normal variable, the value following an assignment to the attribute immediately reflects the new value even though its impact on the ceiling priority of P is postponed until completion of the protected action in which it is executed.

11/2

D.6 Preemptive Abort

- 1 This clause specifies requirements on the immediacy with which an aborted construct is completed.

Dynamic Semantics

- 2 On a system with a single processor, an aborted construct is completed immediately at the first point that is outside the execution of an abort-deferred operation.

Documentation Requirements

- 3 On a multiprocessor, the implementation shall document any conditions that cause the completion of an aborted construct to be delayed later than what is specified for a single processor.

Metrics

- 4 The implementation shall document the following metrics:

- 5 • The execution time, in processor clock cycles, that it takes for an `abort_statement` to cause the completion of the aborted task. This is measured in a situation where a task T2 preempts task T1 and aborts T1. T1 does not have any finalization code. T2 shall verify that T1 has terminated, by means of the `Terminated` attribute.
- 6 • On a multiprocessor, an upper bound in seconds, on the time that the completion of an aborted task can be delayed beyond the point that it is required for a single processor.

- 7/2 • An upper bound on the execution time of an `asynchronous_select`, in processor clock cycles. This is measured between a point immediately before a task T1 executes a protected operation Pr.Set that makes the condition of an `entry_barrier` Pr.Wait True, and the point where task T2 resumes execution immediately after an entry call to Pr.Wait in an `asynchronous_select`. T1 preempts T2 while T2 is executing the abortable part, and then blocks itself so that T2 can execute. The execution time of T1 is measured separately, and subtracted.

- 8 • An upper bound on the execution time of an `asynchronous_select`, in the case that no asynchronous transfer of control takes place. This is measured between a point immediately before a task executes the `asynchronous_select` with a nonnull abortable part, and the point where the task continues execution immediately after it. The execution time of the abortable part is subtracted.

Implementation Advice

- 9 Even though the `abort_statement` is included in the list of potentially blocking operations (see 9.5.1), it is recommended that this statement be implemented in a way that never requires the task executing the `abort_statement` to block.

- 10 On a multi-processor, the delay associated with aborting a task on another processor should be bounded; the implementation should use periodic polling, if necessary, to achieve this.

NOTES

- 11 32 Abortion does not change the active or base priority of the aborted task.
- 12 33 Abortion cannot be more immediate than is allowed by the rules for deferral of abortion during finalization and in protected actions.

D.7 Tasking Restrictions

This clause defines restrictions that can be used with a pragma Restrictions (see 13.12) to facilitate the construction of highly efficient tasking run-time systems.

Static Semantics

The following *restriction_identifiers* are language defined:

No_Task_Hierarchy

All (nonenvironment) tasks depend directly on the environment task of the partition.

No_Nested_Finalization

Objects of a type that needs finalization (see 7.6) and access types that designate a type that needs finalization shall be declared only at library level.

No_Abort_Statements

There are no `abort_statements`, and there are no calls on `Task_Identification.Abort_Task`.

No_Terminate_Alternatives

There are no `selective_accepts` with `terminate_alternatives`.

No_Task_Allocators

There are no `allocators` for task types or types containing task subcomponents.

No_Implicit_Heap_Allocations

There are no operations that implicitly require heap storage allocation to be performed by the implementation. The operations that implicitly require heap storage allocation are implementation defined.

No_Dynamic_Priorities

There are no semantic dependences on the package `Dynamic_Priorities`, and no occurrences of the attribute `Priority`.

No_Dynamic_Attachment

There is no call to any of the operations defined in package `Interrupts` (`Is_Reserved`, `Is_Attached`, `Current_Handler`, `Attach_Handler`, `Exchange_Handler`, `Detach_Handler`, and `Reference`).

No_Local_Protected_Objects

Protected objects shall be declared only at library level.

No_Local_Timing_Events

`Timing_Events` shall be declared only at library level.

No_Protected_Type_Allocators

There are no `allocators` for protected types or types containing protected type subcomponents.

No_Relative_Delay

There are no `delay_relative_statements`.

No_Requeue_Statements

There are no `requeue_statements`.

No_Select_Statements

There are no `select_statements`.

10.7/2 No_Specific_Termination_Handlers

There are no calls to the Set_Specific_Handler and Specific_Handler subprograms in Task_Termination.

10.8/2 Simple_Barriers

The Boolean expression in an entry barrier shall be either a static Boolean expression or a Boolean component of the enclosing protected object.

11 The following *restriction_parameter_identifiers* are language defined:

12 Max_Select_Alternatives

Specifies the maximum number of alternatives in a selective_accept.

13 Max_Task_Entries

Specifies the maximum number of entries per task. The bounds of every entry family of a task unit shall be static, or shall be defined by a discriminant of a subtype whose corresponding bound is static. A value of zero indicates that no rendezvous are possible.

14 Max_ProtectedEntries

Specifies the maximum number of entries per protected type. The bounds of every entry family of a protected unit shall be static, or shall be defined by a discriminant of a subtype whose corresponding bound is static.

Dynamic Semantics

15/2 The following *restriction_identifier* is language defined:

15.1/2 No_Task_Termination

All tasks are non-terminating. It is implementation-defined what happens if a task attempts to terminate. If there is a fall-back handler (see C.7.3) set for the partition it should be called when the first task attempts to terminate.

16 The following *restriction_parameter_identifiers* are language defined:

17/1 Max_Storage_At_Blocking

Specifies the maximum portion (in storage elements) of a task's Storage_Size that can be retained by a blocked task. If an implementation chooses to detect a violation of this restriction, Storage_Error should be raised; otherwise, the behavior is implementation defined.

18/1 Max_Asynchronous_Select_Nesting

Specifies the maximum dynamic nesting level of asynchronous_selects. A value of zero prevents the use of any asynchronous_select and, if a program contains an asynchronous_select, it is illegal. If an implementation chooses to detect a violation of this restriction for values other than zero, Storage_Error should be raised; otherwise, the behavior is implementation defined.

19/1 Max_Tasks Specifies the maximum number of task creations that may be executed over the lifetime of a partition, not counting the creation of the environment task. A value of zero prevents any task creation and, if a program contains a task creation, it is illegal. If an implementation chooses to detect a violation of this restriction, Storage_Error should be raised; otherwise, the behavior is implementation defined.

19.1/2 Max_Entry_Queue_Length

Max_Entry_Queue_Length defines the maximum number of calls that are queued on an entry. Violation of this restriction results in the raising of Program_Error at the point of the call or requeue.

It is implementation defined whether the use of pragma Restrictions results in a reduction in executable program size, storage requirements, or execution time. If possible, the implementation should provide quantitative descriptions of such effects for each restriction.

Implementation Advice

When feasible, the implementation should take advantage of the specified restrictions to produce a more efficient implementation.

NOTES

34 The above Storage_Checks can be suppressed with pragma Suppress.

D.8 Monotonic Time

This clause specifies a high-resolution, monotonic clock package.

Static Semantics

The following language-defined library package exists:

```
package Ada.Real_Time is
    type Time is private;
    Time_First : constant Time;
    Time_Last : constant Time;
    Time_Unit : constant := implementation-defined-real-number;

    type Time_Span is private;
    Time_Span_First : constant Time_Span;
    Time_Span_Last : constant Time_Span;
    Time_Span_Zero : constant Time_Span;
    Time_Span_Unit : constant Time_Span;

    Tick : constant Time_Span;
    function Clock return Time;

    function "+" (Left : Time; Right : Time_Span) return Time;
    function "+" (Left : Time_Span; Right : Time) return Time;
    function "-" (Left : Time; Right : Time_Span) return Time;
    function "-" (Left : Time; Right : Time) return Time_Span;

    function "<" (Left, Right : Time) return Boolean;
    function "<=" (Left, Right : Time) return Boolean;
    function ">" (Left, Right : Time) return Boolean;
    function ">=" (Left, Right : Time) return Boolean;

    function "+" (Left, Right : Time_Span) return Time_Span;
    function "-" (Left, Right : Time_Span) return Time_Span;
    function "-" (Right : Time_Span) return Time_Span;
    function "*" (Left : Time_Span; Right : Integer) return Time_Span;
    function "*" (Left : Integer; Right : Time_Span) return Time_Span;
    function "/" (Left, Right : Time_Span) return Integer;
    function "/" (Left : Time_Span; Right : Integer) return Time_Span;

    function "abs" (Right : Time_Span) return Time_Span;
```

This paragraph was deleted.

```
    function "<" (Left, Right : Time_Span) return Boolean;
    function "<=" (Left, Right : Time_Span) return Boolean;
    function ">" (Left, Right : Time_Span) return Boolean;
    function ">=" (Left, Right : Time_Span) return Boolean;

    function To_Duration (TS : Time_Span) return Duration;
    function To_Time_Span (D : Duration) return Time_Span;
```

```

14/2      function Nanoseconds  (NS : Integer) return Time_Span;
function Microseconds (US : Integer) return Time_Span;
function Milliseconds (MS : Integer) return Time_Span;
function Seconds      (S  : Integer) return Time_Span;
function Minutes       (M  : Integer) return Time_Span;

15      type Seconds_Count is range implementation-defined;
16      procedure Split(T : in Time; SC : out Seconds_Count; TS : out Time_Span);
function Time_Of(SC : Seconds_Count; TS : Time_Span) return Time;

17  private
    . . . -- not specified by the language
end Ada.Real_Time;

```

18 In this Annex, *real time* is defined to be the physical time as observed in the external environment. The type Time is a *time type* as defined by 9.6; values of this type may be used in a delay_until_statement. Values of this type represent segments of an ideal time line. The set of values of the type Time corresponds one-to-one with an implementation-defined range of mathematical integers.

19 The Time value I represents the half-open real time interval that starts with $E+I\cdot Time_Unit$ and is limited by $E+(I+1)\cdot Time_Unit$, where Time_Unit is an implementation-defined real number and E is an unspecified origin point, the *epoch*, that is the same for all values of the type Time. It is not specified by the language whether the time values are synchronized with any standard time reference. For example, E can correspond to the time of system initialization or it can correspond to the epoch of some time standard.

20 Values of the type Time_Span represent length of real time duration. The set of values of this type corresponds one-to-one with an implementation-defined range of mathematical integers. The Time_Span value corresponding to the integer I represents the real-time duration $I\cdot Time_Unit$.

21 Time_First and Time_Last are the smallest and largest values of the Time type, respectively. Similarly, Time_Span_First and Time_Span_Last are the smallest and largest values of the Time_Span type, respectively.

22 A value of type Seconds_Count represents an elapsed time, measured in seconds, since the epoch.

Dynamic Semantics

23 Time_Unit is the smallest amount of real time representable by the Time type; it is expressed in seconds. Time_Span_Unit is the difference between two successive values of the Time type. It is also the smallest positive value of type Time_Span. Time_Unit and Time_Span_Unit represent the same real time duration. A *clock tick* is a real time interval during which the clock value (as observed by calling the Clock function) remains constant. Tick is the average length of such intervals.

24/2 The function To_Duration converts the value TS to a value of type Duration. Similarly, the function To_Time_Span converts the value D to a value of type Time_Span. For To_Duration, the result is rounded to the nearest value of type Duration (away from zero if exactly halfway between two values). If the result is outside the range of Duration, Constraint_Error is raised. For To_Time_Span, the value of D is first rounded to the nearest integral multiple of Time_Unit, away from zero if exactly halfway between two multiples. If the rounded value is outside the range of Time_Span, Constraint_Error is raised. Otherwise, the value is converted to the type Time_Span.

25 To_Duration(Time_Span_Zero) returns 0.0, and To_Time_Span(0.0) returns Time_Span_Zero.

26/2 The functions Nanoseconds, Microseconds, Milliseconds, Seconds, and Minutes convert the input parameter to a value of the type Time_Span. NS, US, MS, S, and M are interpreted as a number of nanoseconds, microseconds, milliseconds, seconds, and minutes respectively. The input parameter is first converted to seconds and rounded to the nearest integral multiple of Time_Unit, away from zero if exactly

halfway between two multiples. If the rounded value is outside the range of Time_Span, Constraint_Error is raised. Otherwise, the rounded value is converted to the type Time_Span.

The effects of the operators on Time and Time_Span are as for the operators defined for integer types. 27

The function Clock returns the amount of time since the epoch. 28

The effects of the Split and Time_Of operations are defined as follows, treating values of type Time, Time_Span, and Seconds_Count as mathematical integers. The effect of Split(T,SC,TS) is to set SC and TS to values such that $T * \text{Time_Unit} = SC * 1.0 + TS * \text{Time_Unit}$, and $0.0 \leq TS * \text{Time_Unit} < 1.0$. The value returned by Time_Of(SC,TS) is the value T such that $T * \text{Time_Unit} = SC * 1.0 + TS * \text{Time_Unit}$. 29

Implementation Requirements

The range of Time values shall be sufficient to uniquely represent the range of real times from program start-up to 50 years later. Tick shall be no greater than 1 millisecond. Time_Unit shall be less than or equal to 20 microseconds. 30

Time_Span_First shall be no greater than -3600 seconds, and Time_Span_Last shall be no less than 3600 seconds. 31

A *clock jump* is the difference between two successive distinct values of the clock (as observed by calling the Clock function). There shall be no backward clock jumps. 32

Documentation Requirements

The implementation shall document the values of Time_First, Time_Last, Time_Span_First, Time_Span_Last, Time_Span_Unit, and Tick. 33

The implementation shall document the properties of the underlying time base used for the clock and for type Time, such as the range of values supported and any relevant aspects of the underlying hardware or operating system facilities used. 34

The implementation shall document whether or not there is any synchronization with external time references, and if such synchronization exists, the sources of synchronization information, the frequency of synchronization, and the synchronization method applied. 35

The implementation shall document any aspects of the external environment that could interfere with the clock behavior as defined in this clause. 36/1

Metrics

For the purpose of the metrics defined in this clause, real time is defined to be the International Atomic Time (TAI). 37

The implementation shall document the following metrics: 38

- An upper bound on the real-time duration of a clock tick. This is a value D such that if t1 and t2 are any real times such that $t1 < t2$ and $\text{Clock}_{t1} = \text{Clock}_{t2}$, then $t2 - t1 \leq D$. 39
- An upper bound on the size of a clock jump. 40
- An upper bound on the *drift rate* of Clock with respect to real time. This is a real number D such that 41

$$E * (1 - D) \leq (\text{Clock}_{t+E} - \text{Clock}_t) \leq E * (1 + D)$$

provided that: $\text{Clock}_t + E * (1 + D) \leq \text{Time_Last}$. 42

- 43 • where Clock_t is the value of Clock at time t , and E is a real time duration not less than 24 hours.
The value of E used for this metric shall be reported.
- 44 • An upper bound on the execution time of a call to the Clock function, in processor clock cycles.
- 45 • Upper bounds on the execution times of the operators of the types Time and Time_Span, in
processor clock cycles.

Implementation Permissions

46 Implementations targeted to machines with word size smaller than 32 bits need not support the full range
and granularity of the Time and Time_Span types.

Implementation Advice

- 47 When appropriate, implementations should provide configuration mechanisms to change the value of Tick.
- 48 It is recommended that Calendar.Clock and Real_Time.Clock be implemented as transformations of the
same time base.
- 49 It is recommended that the “best” time base which exists in the underlying system be available to the
application through Clock. “Best” may mean highest accuracy or largest range.

NOTES

50 35 The rules in this clause do not imply that the implementation can protect the user from operator or installation errors
which could result in the clock being set incorrectly.

51 36 Time_Unit is the granularity of the Time type. In contrast, Tick represents the granularity of Real_Time.Clock. There
is no requirement that these be the same.

D.9 Delay Accuracy

- 1 This clause specifies performance requirements for the `delay_statement`. The rules apply both to `delay_-relative_statement` and to `delay_until_statement`. Similarly, they apply equally to a simple `delay_-statement` and to one which appears in a `delay_alternative`.

Dynamic Semantics

2 The effect of the `delay_statement` for Real_Time.Time is defined in terms of Real_Time.Clock:

- 3 • If C_1 is a value of Clock read before a task executes a `delay_relative_statement` with duration D ,
and C_2 is a value of Clock read after the task resumes execution following that `delay_statement`,
then $C_2 - C_1 \geq D$.
- 4 • If C is a value of Clock read after a task resumes execution following a `delay_until_statement`
with Real_Time.Time value T , then $C \geq T$.

5 A simple `delay_statement` with a negative or zero value for the expiration time does not cause the calling
task to be blocked; it is nevertheless a potentially blocking operation (see 9.5.1).

62 When a `delay_statement` appears in a `delay_alternative` of a `timed_entry_call` the selection of the entry
call is attempted, regardless of the specified expiration time. When a `delay_statement` appears in a
`select_alternative`, and a call is queued on one of the open entries, the selection of that entry call proceeds,
regardless of the value of the delay expression.

Documentation Requirements

- 7 The implementation shall document the minimum value of the delay expression of a
`delay_relative_statement` that causes the task to actually be blocked.

The implementation shall document the minimum difference between the value of the delay expression of a `delay_until_statement` and the value of `Real_Time.Clock`, that causes the task to actually be blocked.

Metrics

The implementation shall document the following metrics:

- An upper bound on the execution time, in processor clock cycles, of a `delay_relative_statement` whose requested value of the delay expression is less than or equal to zero.
- An upper bound on the execution time, in processor clock cycles, of a `delay_until_statement` whose requested value of the delay expression is less than or equal to the value of `Real_Time.Clock` at the time of executing the statement. Similarly, for `Calendar.Clock`.
- An upper bound on the *lateness* of a `delay_relative_statement`, for a positive value of the delay expression, in a situation where the task has sufficient priority to preempt the processor as soon as it becomes ready, and does not need to wait for any other execution resources. The upper bound is expressed as a function of the value of the delay expression. The lateness is obtained by subtracting the value of the delay expression from the *actual duration*. The actual duration is measured from a point immediately before a task executes the `delay_statement` to a point immediately after the task resumes execution following this statement.
- An upper bound on the lateness of a `delay_until_statement`, in a situation where the value of the requested expiration time is after the time the task begins executing the statement, the task has sufficient priority to preempt the processor as soon as it becomes ready, and it does not need to wait for any other execution resources. The upper bound is expressed as a function of the difference between the requested expiration time and the clock value at the time the statement begins execution. The lateness of a `delay_until_statement` is obtained by subtracting the requested expiration time from the real time that the task resumes execution following this statement.

NOTES

This paragraph was deleted.

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D.10 Synchronous Task Control

This clause describes a language-defined private semaphore (suspension object), which can be used for *two-stage suspend* operations and as a simple building block for implementing higher-level queues.

Static Semantics

The following language-defined package exists:

```
package Ada.Synchronous_Task_Control is
  pragma Preelaborate(Synchronous_Task_Control);
  type Suspension_Object is limited private;
  procedure Set_True(S : in out Suspension_Object);
  procedure Set_False(S : in out Suspension_Object);
  function Current_State(S : Suspension_Object) return Boolean;
  procedure Suspend_Until_True(S : in out Suspension_Object);
  private
    ... -- not specified by the language
end Ada.Synchronous_Task_Control;
```

2

3/2

4

The type `Suspension_Object` is a by-reference type.

Dynamic Semantics

An object of the type `Suspension_Object` has two visible states: True and False. Upon initialization, its value is set to False.

5

6/2

- 7/2 The operations Set_True and Set_False are atomic with respect to each other and with respect to Suspend_Until_True; they set the state to True and False respectively.
- 8 Current_State returns the current state of the object.
- 9/2 The procedure Suspend_Until_True blocks the calling task until the state of the object S is True; at that point the task becomes ready and the state of the object becomes False.
- 10 Program_Error is raised upon calling Suspend_Until_True if another task is already waiting on that suspension object. Suspend_Until_True is a potentially blocking operation (see 9.5.1).

Implementation Requirements

- 11 The implementation is required to allow the calling of Set_False and Set_True during any protected action, even one that has its ceiling priority in the Interrupt_Priority range.

D.11 Asynchronous Task Control

- 1 This clause introduces a language-defined package to do asynchronous suspend/resume on tasks. It uses a conceptual *held priority* value to represent the task's *held* state.

Static Semantics

- 2 The following language-defined library package exists:

```
3/2   with Ada.Task_Identification;
      package Ada.Asynchronous_Task_Control is
        pragma Preelaborate(Asynchronous_Task_Control);
        procedure Hold(T : in Ada.Task_Identification.Task_Id);
        procedure Continue(T : in Ada.Task_Identification.Task_Id);
        function Is_Held(T : Ada.Task_Identification.Task_Id)
          return Boolean;
      end Ada.Asynchronous_Task_Control;
```

Dynamic Semantics

- 4/2 After the Hold operation has been applied to a task, the task becomes *held*. For each processor there is a conceptual *idle task*, which is always ready. The base priority of the idle task is below System.Any_Priority'First. The *held priority* is a constant of the type Integer whose value is below the base priority of the idle task.
- 4.1/2 For any priority below System.Any_Priority'First, the task dispatching policy is FIFO_Within_Priorities.
- 5/2 The Hold operation sets the state of T to held. For a held task, the active priority is reevaluated as if the base priority of the task were the held priority.
- 6/2 The Continue operation resets the state of T to not-held; its active priority is then reevaluated as determined by the task dispatching policy associated with its base priority.
- 7 The Is_Held function returns True if and only if T is in the held state.
- 8 As part of these operations, a check is made that the task identified by T is not terminated. Tasking_Error is raised if the check fails. Program_Error is raised if the value of T is Null_Task_Id.

Erroneous Execution

- 9 If any operation in this package is called with a parameter T that specifies a task object that no longer exists, the execution of the program is erroneous.

Implementation Permissions

An implementation need not support Asynchronous_Task_Control if it is infeasible to support it in the target environment. 10

NOTES

- 37 It is a consequence of the priority rules that held tasks cannot be dispatched on any processor in a partition (unless they are inheriting priorities) since their priorities are defined to be below the priority of any idle task. 11
- 38 The effect of calling Get_Priority and Set_Priority on a Held task is the same as on any other task. 12
- 39 Calling Hold on a held task or Continue on a non-held task has no effect. 13
- 40 The rules affecting queuing are derived from the above rules, in addition to the normal priority rules: 14
 - When a held task is on the ready queue, its priority is so low as to never reach the top of the queue as long as there are other tasks on that queue. 15
 - If a task is executing in a protected action, inside a rendezvous, or is inheriting priorities from other sources (e.g. when activated), it continues to execute until it is no longer executing the corresponding construct. 16
 - If a task becomes held while waiting (as a caller) for a rendezvous to complete, the active priority of the accepting task is not affected. 17
 - If a task becomes held while waiting in a selective_accept, and an entry call is issued to one of the open entries, the corresponding accept_alternative executes. When the rendezvous completes, the active priority of the accepting task is lowered to the held priority (unless it is still inheriting from other sources), and the task does not execute until another Continue. 18/1
 - The same holds if the held task is the only task on a protected entry queue whose barrier becomes open. The corresponding entry body executes. 19

D.12 Other Optimizations and Determinism Rules

This clause describes various requirements for improving the response and determinism in a real-time system. 1

Implementation Requirements

If the implementation blocks interrupts (see C.3) not as a result of direct user action (e.g. an execution of a protected action) there shall be an upper bound on the duration of this blocking. 2

The implementation shall recognize entry-less protected types. The overhead of acquiring the execution resource of an object of such a type (see 9.5.1) shall be minimized. In particular, there should not be any overhead due to evaluating entry_barrier conditions. 3

Unchecked_Deallocation shall be supported for terminated tasks that are designated by access types, and shall have the effect of releasing all the storage associated with the task. This includes any run-time system or heap storage that has been implicitly allocated for the task by the implementation. 4

Documentation Requirements

The implementation shall document the upper bound on the duration of interrupt blocking caused by the implementation. If this is different for different interrupts or interrupt priority levels, it should be documented for each case. 5

Metrics

The implementation shall document the following metric: 6

- The overhead associated with obtaining a mutual-exclusive access to an entry-less protected object. This shall be measured in the following way: 7

- 8 For a protected object of the form:

```

9   protected Lock is
10    procedure Set;
11     function Read return Boolean;
12   private
13     Flag : Boolean := False;
14   end Lock;

15   protected body Lock is
16    procedure Set is
17     begin
18       Flag := True;
19     end Set;
20     function Read return Boolean
21     Begin
22       return Flag;
23     end Read;
24   end Lock;
```

- 11 The execution time, in processor clock cycles, of a call to Set. This shall be measured between the point just before issuing the call, and the point just after the call completes. The function Read shall be called later to verify that Set was indeed called (and not optimized away). The calling task shall have sufficiently high priority as to not be preempted during the measurement period. The protected object shall have sufficiently high ceiling priority to allow the task to call Set.

- 12 For a multiprocessor, if supported, the metric shall be reported for the case where no contention (on the execution resource) exists from tasks executing on other processors.

D.13 Run-time Profiles

- 1/2 This clause specifies a mechanism for defining run-time profiles.

Syntax

- 2/2 The form of a **pragma Profile** is as follows:

```
pragma Profile (profile_identifier {, profile_pragma_argument_association});
```

Legality Rules

- 4/2 The *profile_identifier* shall be the name of a run-time profile. The semantics of any *profile_pragma_argument_associations* are defined by the run-time profile specified by the *profile_identifier*.

Static Semantics

- 5/2 A profile is equivalent to the set of configuration pragmas that is defined for each run-time profile.

Post-Compilation Rules

- 6/2 A **pragma Profile** is a configuration pragma. There may be more than one **pragma Profile** for a partition.

D.13.1 The Ravenscar Profile

This clause defines the Ravenscar profile.

1/2

Legality Rules

The *profile_identifier* Ravenscar is a run-time profile. For run-time profile Ravenscar, there shall be no *profile_pragma_argument_associations*.

2/2

Static Semantics

The run-time profile Ravenscar is equivalent to the following set of pragmas:

3/2

```
pragma Task_Dispatching_Policy (FIFO_Within_Priorities);
pragma Locking_Policy (Ceiling_Locking);
pragma Detect_Blocking;
pragma Restrictions (
    No_Abort_Statements,
    No_Dynamic_Attachment,
    No_Dynamic_Priorities,
    No_Implicit_Heap_Allocations,
    No_Local_Protected_Objects,
    No_Local_Timing_Events,
    No_Protected_Type_Allocators,
    No_Relative_Delay,
    No_Requeue_Statements,
    No_Select_Statements,
    No_Specific_Termination_Handlers,
    No_Task_Allocators,
    No_Task_Hierarchy,
    No_Task_Termination,
    Simple_Barriers,
    Max_Entry_Queue_Length => 1,
    Max_ProtectedEntries => 1,
    Max_TaskEntries => 0,
    No_Dependence => Ada.Asynchronous_Task_Control,
    No_Dependence => Ada.Calendar,
    No_Dependence => Ada.Execution_Time.Group_Budget,
    No_Dependence => Ada.Execution_Time.Timers,
    No_Dependence => Ada.Task_Attributes);
```

4/2

NOTES

41 The effect of the *Max_Entry_Queue_Length => 1* restriction applies only to protected entry queues due to the accompanying restriction of *Max_TaskEntries => 0*.

5/2

D.14 Execution Time

- 1/2 This clause describes a language-defined package to measure execution time.

Static Semantics

- 2/2 The following language-defined library package exists:

```

3/2   with Ada.Task_Identification;
      with Ada.Real_Time; use Ada.Real_Time;
      package Ada.Execution_Time is
4/2     type CPU_Time is private;
        CPU_Time_First : constant CPU_Time;
        CPU_Time_Last  : constant CPU_Time;
        CPU_Time_Unit  : constant := implementation-defined-real-number;
        CPU_Tick       : constant Time_Span;
5/2     function Clock
              (T : Ada.Task_Identification.Task_Id
               := Ada.Task_Identification.Current_Task)
             return CPU_Time;
6/2     function "+"  (Left : CPU_Time; Right : Time_Span) return CPU_Time;
     function "+"  (Left : Time_Span; Right : CPU_Time) return CPU_Time;
     function "-"  (Left : CPU_Time; Right : Time_Span) return CPU_Time;
     function "-"  (Left : CPU_Time; Right : CPU_Time) return Time_Span;
7/2     function "<"  (Left, Right : CPU_Time) return Boolean;
     function "<=" (Left, Right : CPU_Time) return Boolean;
     function ">"  (Left, Right : CPU_Time) return Boolean;
     function ">=" (Left, Right : CPU_Time) return Boolean;
8/2     procedure Split
              (T : in CPU_Time; SC : out Seconds_Count; TS : out Time_Span);
9/2     function Time_Of (SC : Seconds_Count;
                         TS : Time_Span := Time_Span_Zero) return CPU_Time;
10/2    private
           ... -- not specified by the language
      end Ada.Execution_Time;
```

- 11/2 The *execution time* or CPU time of a given task is defined as the time spent by the system executing that task, including the time spent executing run-time or system services on its behalf. The mechanism used to measure execution time is implementation defined. It is implementation defined which task, if any, is charged the execution time that is consumed by interrupt handlers and run-time services on behalf of the system.

- 12/2 The type CPU_Time represents the execution time of a task. The set of values of this type corresponds one-to-one with an implementation-defined range of mathematical integers.

- 13/2 The CPU_Time value I represents the half-open execution-time interval that starts with I*CPU_Time_Unit and is limited by (I+1)*CPU_Time_Unit, where CPU_Time_Unit is an implementation-defined real number. For each task, the execution time value is set to zero at the creation of the task.

- 14/2 CPU_Time_First and CPU_Time_Last are the smallest and largest values of the CPU_Time type, respectively.

Dynamic Semantics

- 15/2 CPU_Time_Unit is the smallest amount of execution time representable by the CPU_Time type; it is expressed in seconds. A *CPU clock tick* is an execution time interval during which the clock value (as

observed by calling the `Clock` function) remains constant. `CPU_Tick` is the average length of such intervals.

The effects of the operators on `CPU_Time` and `Time_Span` are as for the operators defined for integer types. 16/2

The function `Clock` returns the current execution time of the task identified by `T`; `Tasking_Error` is raised if that task has terminated; `Program_Error` is raised if the value of `T` is `Task_Identification.Null_Task_Id`. 17/2

The effects of the `Split` and `Time_Of` operations are defined as follows, treating values of type `CPU_Time`, `Time_Span`, and `Seconds_Count` as mathematical integers. The effect of `Split` (`T, SC, TS`) is to set `SC` and `TS` to values such that $T * \text{CPU_Time_Unit} = SC * 1.0 + TS * \text{CPU_Time_Unit}$, and $0.0 \leq TS * \text{CPU_Time_Unit} < 1.0$. The value returned by `Time_Of(SC, TS)` is the execution-time value `T` such that $T * \text{CPU_Time_Unit} = SC * 1.0 + TS * \text{CPU_Time_Unit}$. 18/2

Erroneous Execution

For a call of `Clock`, if the task identified by `T` no longer exists, the execution of the program is erroneous. 19/2

Implementation Requirements

The range of `CPU_Time` values shall be sufficient to uniquely represent the range of execution times from the task start-up to 50 years of execution time later. `CPU_Tick` shall be no greater than 1 millisecond. 20/2

Documentation Requirements

The implementation shall document the values of `CPU_Time_First`, `CPU_Time_Last`, `CPU_Time_Unit`, and `CPU_Tick`. 21/2

The implementation shall document the properties of the underlying mechanism used to measure execution times, such as the range of values supported and any relevant aspects of the underlying hardware or operating system facilities used. 22/2

Metrics

The implementation shall document the following metrics: 23/2

- An upper bound on the execution-time duration of a clock tick. This is a value `D` such that if `t1` and `t2` are any execution times of a given task such that $t1 < t2$ and $\text{Clock}_{t1} = \text{Clock}_{t2}$ then $t2 - t1 \leq D$. 24/2
- An upper bound on the size of a clock jump. A clock jump is the difference between two successive distinct values of an execution-time clock (as observed by calling the `Clock` function with the same `Task_Id`). 25/2
- An upper bound on the execution time of a call to the `Clock` function, in processor clock cycles. 26/2
- Upper bounds on the execution times of the operators of the type `CPU_Time`, in processor clock cycles. 27/2

Implementation Permissions

Implementations targeted to machines with word size smaller than 32 bits need not support the full range and granularity of the `CPU_Time` type. 28/2

Implementation Advice

When appropriate, implementations should provide configuration mechanisms to change the value of `CPU_Tick`. 29/2

D.14.1 Execution Time Timers

- 1/2 This clause describes a language-defined package that provides a facility for calling a handler when a task has used a defined amount of CPU time.

Static Semantics

- 2/2 The following language-defined library package exists:

```

3/2   with System;
4/2   package Ada.Execution_Time.Timers is
5/2     type Timer (T : not null access constant
                   Ada.Task_Identification.Task_Id) is
6/2       tagged limited private;
7/2     type Timer_Handler is
9/2       access protected procedure (TM : in out Timer);
10/2    Min_Handler_Ceiling : constant System.Any_Priority := implementation-defined;
11/2    procedure Set_Handler (TM      : in out Timer;
12/2                           In_Time : in Time_Span;
13/2                           Handler : in Timer_Handler);
14/2    procedure Set_Handler (TM      : in out Timer;
15/2                           At_Time : in CPU_Time;
16/2                           Handler : in Timer_Handler);
17/2    function Current_Handler (TM : Timer) return Timer_Handler;
18/2    procedure Cancel_Handler (TM      : in out Timer;
19/2                               Cancelled : out Boolean);
20/2    function Time_Remaining (TM : Timer) return Time_Span;
21/2    Timer_Resource_Error : exception;
22/2  private
23/2    ... -- not specified by the language
24/2 end Ada.Execution_Time.Timers;
```

- 11/2 The type `Timer` represents an execution-time event for a single task and is capable of detecting execution-time overruns. The access discriminant `T` identifies the task concerned. The type `Timer` needs finalization (see 7.6).

- 12/2 An object of type `Timer` is said to be *set* if it is associated with a non-null value of type `Timer_Handler` and *cleared* otherwise. All `Timer` objects are initially cleared.

- 13/2 The type `Timer_Handler` identifies a protected procedure to be executed by the implementation when the timer expires. Such a protected procedure is called a *handler*.

Dynamic Semantics

- 14/2 When a `Timer` object is created, or upon the first call of a `Set_Handler` procedure with the timer as parameter, the resources required to operate an execution-time timer based on the associated execution-time clock are allocated and initialized. If this operation would exceed the available resources, `Timer_Resource_Error` is raised.

- 15/2 The procedures `Set_Handler` associate the handler `Handler` with the timer `TM`; if `Handler` is `null`, the timer is cleared, otherwise it is set. The first procedure `Set_Handler` loads the timer `TM` with an interval specified by the `Time_Span` parameter. In this mode, the timer `TM` *expires* when the execution time of the task identified by `TM.T.all` has increased by `In_Time`; if `In_Time` is less than or equal to zero, the timer expires immediately. The second procedure `Set_Handler` loads the timer `TM` with the absolute value specified by `At_Time`. In this mode, the timer `TM` expires when the execution time of the task identified

by `TM.T.all` reaches `At_Time`; if the value of `At_Time` has already been reached when `Set_Handler` is called, the timer expires immediately.

A call of a procedure `Set_Handler` for a timer that is already set replaces the handler and the (absolute or relative) execution time; if `Handler` is not `null`, the timer remains set. 16/2

When a timer expires, the associated handler is executed, passing the timer as parameter. The initial action of the execution of the handler is to clear the event. 17/2

The function `Current_Handler` returns the handler associated with the timer `TM` if that timer is set; otherwise it returns `null`. 18/2

The procedure `Cancel_Handler` clears the timer if it is set. `Cancelled` is assigned `True` if the timer was set prior to it being cleared; otherwise it is assigned `False`. 19/2

The function `Time_Remaining` returns the execution time interval that remains until the timer `TM` would expire, if that timer is set; otherwise it returns `Time_Span_Zero`. 20/2

The constant `Min_Handler_Ceiling` is the minimum ceiling priority required for a protected object with a handler to ensure that no ceiling violation will occur when that handler is invoked. 21/2

As part of the finalization of an object of type `Timer`, the timer is cleared. 22/2

For all the subprograms defined in this package, `Tasking_Error` is raised if the task identified by `TM.T.all` has terminated, and `Program_Error` is raised if the value of `TM.T.all` is `Task_Identification.Null_Task_Id`. 23/2

An exception propagated from a handler invoked as part of the expiration of a timer has no effect. 24/2

Erroneous Execution

For a call of any of the subprograms defined in this package, if the task identified by `TM.T.all` no longer exists, the execution of the program is erroneous. 25/2

Implementation Requirements

For a given `Timer` object, the implementation shall perform the operations declared in this package atomically with respect to any of these operations on the same `Timer` object. The replacement of a handler by a call of `Set_Handler` shall be performed atomically with respect to the execution of the handler. 26/2

When an object of type `Timer` is finalized, the system resources used by the timer shall be deallocated. 27/2

Implementation Permissions

Implementations may limit the number of timers that can be defined for each task. If this limit is exceeded then `Timer_Resource_Error` is raised. 28/2

NOTES

42 A `Timer_Handler` can be associated with several `Timer` objects. 29/2

D.14.2 Group Execution Time Budgets

- 1/2 This clause describes a language-defined package to assign execution time budgets to groups of tasks.

Static Semantics

- 2/2 The following language-defined library package exists:

```

3/2   with System;
4/2   package Ada.Execution_Time.Group_Budgets is
5/2     type Group_Budget is tagged limited private;
6/2     type Group_Budget_Handler is access
9/2       protected procedure (GB : in out Group_Budget);
7/2     type Task_Array is array (Positive range <>) of
10/2      Ada.Task_Identification.Task_Id;
11/2      Min_Handler_Ceiling : constant System.Any_Priority := 
12/2        implementation-defined;
13/2      procedure Add_Task (GB : in out Group_Budget;
14/2        T : in Ada.Task_Identification.Task_Id);
15/2      procedure Remove_Task (GB: in out Group_Budget;
16/2        T : in Ada.Task_Identification.Task_Id);
17/2      function Is_Member (GB : Group_Budget;
18/2        T : Ada.Task_Identification.Task_Id) return Boolean;
19/2      function Is_A_Group_Member
20/2        (T : Ada.Task_Identification.Task_Id) return Boolean;
21/2      function Members (GB : Group_Budget) return Task_Array;
22/2      procedure Replenish (GB : in out Group_Budget; To : in Time_Span);
23/2      procedure Add (GB : in out Group_Budget; Interval : in Time_Span);
24/2      function Budget_Has_Expired (GB : Group_Budget) return Boolean;
25/2      function Budget_Remaining (GB : Group_Budget) return Time_Span;
26/2      procedure Set_Handler (GB          : in out Group_Budget;
27/2                                Handler : in Group_Budget_Handler);
28/2      function Current_Handler (GB : Group_Budget)
29/2        return Group_Budget_Handler;
30/2      procedure Cancel_Handler (GB          : in out Group_Budget;
31/2                                Cancelled : out Boolean);
32/2      Group_Budget_Error : exception;
33/2  private
34/2    --  not specified by the language
35/2 end Ada.Execution_Time.Group_Budgets;
```

- 13/2 The type `Group_Budget` represents an execution time budget to be used by a group of tasks. The type `Group_Budget` needs finalization (see 7.6). A task can belong to at most one group. Tasks of any priority can be added to a group.

- 14/2 An object of type `Group_Budget` has an associated nonnegative value of type `Time_Span` known as its *budget*, which is initially `Time_Span_Zero`. The type `Group_Budget_Handler` identifies a protected procedure to be executed by the implementation when the budget is *exhausted*, that is, reaches zero. Such a protected procedure is called a *handler*.

- 15/2 An object of type `Group_Budget` also includes a handler, which is a value of type `Group_Budget_Handler`. The handler of the object is said to be *set* if it is not null and *cleared* otherwise. The handler of all `Group_Budget` objects is initially cleared.

Dynamic Semantics

- 16/2 The procedure `Add_Task` adds the task identified by `T` to the group `GB`; if that task is already a member of some other group, `Group_Budget_Error` is raised.

The procedure Remove_Task removes the task identified by T from the group GB; if that task is not a member of the group GB, Group_Budget_Error is raised. After successful execution of this procedure, the task is no longer a member of any group.	17/2
The function Is_Member returns True if the task identified by T is a member of the group GB; otherwise it return False.	18/2
The function Is_A_Group_Member returns True if the task identified by T is a member of some group; otherwise it returns False.	19/2
The function Members returns an array of values of type Task_Identification.Task_Id identifying the members of the group GB. The order of the components of the array is unspecified.	20/2
The procedure Replenish loads the group budget GB with To as the Time_Span value. The exception Group_Budget_Error is raised if the Time_Span value To is non-positive. Any execution of any member of the group of tasks results in the budget counting down, unless exhausted. When the budget becomes exhausted (reaches Time_Span_Zero), the associated handler is executed if the handler of group budget GB is set. Nevertheless, the tasks continue to execute.	21/2
The procedure Add modifies the budget of the group GB. A positive value for Interval increases the budget. A negative value for Interval reduces the budget, but never below Time_Span_Zero. A zero value for Interval has no effect. A call of procedure Add that results in the value of the budget going to Time_Span_Zero causes the associated handler to be executed if the handler of the group budget GB is set.	22/2
The function Budget_Has_Expired returns True if the budget of group GB is exhausted (equal to Time_Span_Zero); otherwise it returns False.	23/2
The function Budget_Remaining returns the remaining budget for the group GB. If the budget is exhausted it returns Time_Span_Zero. This is the minimum value for a budget.	24/2
The procedure Set_Handler associates the handler Handler with the Group_Budget GB; if Handler is null , the handler of Group_Budget is cleared, otherwise it is set.	25/2
A call of Set_Handler for a Group_Budget that already has a handler set replaces the handler; if Handler is not null , the handler for Group_Budget remains set.	26/2
The function Current_Handler returns the handler associated with the group budget GB if the handler for that group budget is set; otherwise it returns null .	27/2
The procedure Cancel_Handler clears the handler for the group budget if it is set. Cancelled is assigned True if the handler for the group budget was set prior to it being cleared; otherwise it is assigned False.	28/2
The constant Min_Handler_Ceiling is the minimum ceiling priority required for a protected object with a handler to ensure that no ceiling violation will occur when that handler is invoked.	29/2
The precision of the accounting of task execution time to a Group_Budget is the same as that defined for execution-time clocks from the parent package.	30/2
As part of the finalization of an object of type Group_Budget all member tasks are removed from the group identified by that object.	31/2
If a task is a member of a Group_Budget when it terminates then as part of the finalization of the task it is removed from the group.	32/2
For all the operations defined in this package, Tasking_Error is raised if the task identified by T has terminated, and Program_Error is raised if the value of T is Task_Identification.Null_Task_Id.	33/2

- 34/2 An exception propagated from a handler invoked when the budget of a group of tasks becomes exhausted has no effect.

Erroneous Execution

- 35/2 For a call of any of the subprograms defined in this package, if the task identified by T no longer exists, the execution of the program is erroneous.

Implementation Requirements

- 36/2 For a given Group_Budget object, the implementation shall perform the operations declared in this package atomically with respect to any of these operations on the same Group_Budget object. The replacement of a handler, by a call of Set_Handler, shall be performed atomically with respect to the execution of the handler.

NOTES

- 37/2 43 Clearing or setting of the handler of a group budget does not change the current value of the budget. Exhaustion or loading of a budget does not change whether the handler of the group budget is set or cleared.
- 38/2 44 A Group_Budget_Handler can be associated with several Group_Budget objects.

D.15 Timing Events

- 1/2 This clause describes a language-defined package to allow user-defined protected procedures to be executed at a specified time without the need for a task or a delay statement.

Static Semantics

- 2/2 The following language-defined library package exists:

```

package Ada.Real_Time.Timing_Events is
    type Timing_Event is tagged limited private;
    type Timing_Event_Handler
        is access protected procedure (Event : in out Timing_Event);
    procedure Set_Handler (Event      : in out Timing_Event;
                           At_Time   : in Time;
                           Handler   : in Timing_Event_Handler);
    procedure Set_Handler (Event      : in out Timing_Event;
                           In_Time   : in Time_Span;
                           Handler   : in Timing_Event_Handler);
    function Current_Handler (Event : Timing_Event)
        return Timing_Event_Handler;
    procedure Cancel_Handler (Event      : in out Timing_Event;
                             Cancelled : out Boolean);
    function Time_Of_Event (Event : Timing_Event) return Time;
private
    ... -- not specified by the language
end Ada.Real_Time.Timing_Events;
```

- 8/2 The type Timing_Event represents a time in the future when an event is to occur. The type Timing_Event needs finalization (see 7.6).

- 9/2 An object of type Timing_Event is said to be *set* if it is associated with a non-null value of type Timing_Event_Handler and *cleared* otherwise. All Timing_Event objects are initially cleared.

- 10/2 The type Timing_Event_Handler identifies a protected procedure to be executed by the implementation when the timing event occurs. Such a protected procedure is called a *handler*.

Dynamic Semantics

The procedures `Set_Handler` associate the handler `Handler` with the event `Event`; if `Handler` is `null`, the event is cleared, otherwise it is set. The first procedure `Set_Handler` sets the execution time for the event to be `At_Time`. The second procedure `Set_Handler` sets the execution time for the event to be `Real_Time.Clock + In_Time`.

A call of a procedure `Set_Handler` for an event that is already set replaces the handler and the time of execution; if `Handler` is not `null`, the event remains set.

As soon as possible after the time set for the event, the handler is executed, passing the event as parameter. The handler is only executed if the timing event is in the set state at the time of execution. The initial action of the execution of the handler is to clear the event.

If the `Ceiling_Locking` policy (see D.3) is in effect when a procedure `Set_Handler` is called, a check is made that the ceiling priority of `Handler.all` is `Interrupt_Priority'Last`. If the check fails, `Program_Error` is raised.

If a procedure `Set_Handler` is called with zero or negative `In_Time` or with `At_Time` indicating a time in the past then the handler is executed immediately by the task executing the call of `Set_Handler`. The timing event `Event` is cleared.

The function `Current_Handler` returns the handler associated with the event `Event` if that event is set; otherwise it returns `null`.

The procedure `Cancel_Handler` clears the event if it is set. `Cancelled` is assigned True if the event was set prior to it being cleared; otherwise it is assigned False.

The function `Time_Of_Event` returns the time of the event if the event is set; otherwise it returns `Real_Time.Time_First`.

As part of the finalization of an object of type `Timing_Event`, the `Timing_Event` is cleared.

If several timing events are set for the same time, they are executed in FIFO order of being set.

An exception propagated from a handler invoked by a timing event has no effect.

Implementation Requirements

For a given `Timing_Event` object, the implementation shall perform the operations declared in this package atomically with respect to any of these operations on the same `Timing_Event` object. The replacement of a handler by a call of `Set_Handler` shall be performed atomically with respect to the execution of the handler.

Metrics

The implementation shall document the following metric:

- An upper bound on the lateness of the execution of a handler. That is, the maximum time between when a handler is actually executed and the time specified when the event was set.

Implementation Advice

The protected handler procedure should be executed directly by the real-time clock interrupt mechanism.

NOTES

45 Since a call of `Set_Handler` is not a potentially blocking operation, it can be called from within a handler.

46 A `Timing_Event_Handler` can be associated with several `Timing_Event` objects.

Annex E (normative) Distributed Systems

This Annex defines facilities for supporting the implementation of distributed systems using multiple partitions working cooperatively as part of a single Ada program.

Post-Compilation Rules

A *distributed system* is an interconnection of one or more *processing nodes* (a system resource that has both computational and storage capabilities), and zero or more *storage nodes* (a system resource that has only storage capabilities, with the storage addressable by one or more processing nodes).

A *distributed program* comprises one or more partitions that execute independently (except when they communicate) in a distributed system.

The process of mapping the partitions of a program to the nodes in a distributed system is called *configuring the partitions of the program*.

Implementation Requirements

The implementation shall provide means for explicitly assigning library units to a partition and for the configuring and execution of a program consisting of multiple partitions on a distributed system; the means are implementation defined.

Implementation Permissions

An implementation may require that the set of processing nodes of a distributed system be homogeneous.

NOTES

1 The partitions comprising a program may be executed on differently configured distributed systems or on a non-distributed system without requiring recompilation. A distributed program may be partitioned differently from the same set of library units without recompilation. The resulting execution is semantically equivalent.

2 A distributed program retains the same type safety as the equivalent single partition program.

E.1 Partitions

The partitions of a distributed program are classified as either active or passive.

Post-Compilation Rules

An *active partition* is a partition as defined in 10.2. A *passive partition* is a partition that has no thread of control of its own, whose library units are all preelaborated, and whose data and subprograms are accessible to one or more active partitions.

A passive partition shall include only `library_items` that either are declared pure or are shared passive (see 10.2.1 and E.2.1).

An active partition shall be configured on a processing node. A passive partition shall be configured either on a storage node or on a processing node.

The configuration of the partitions of a program onto a distributed system shall be consistent with the possibility for data references or calls between the partitions implied by their semantic dependences. Any reference to data or call of a subprogram across partitions is called a *remote access*.

Dynamic Semantics

- 6 A library_item is elaborated as part of the elaboration of each partition that includes it. If a normal library unit (see E.2) has state, then a separate copy of the state exists in each active partition that elaborates it. The state evolves independently in each such partition.
- 7 An active partition *terminates* when its environment task terminates. A partition becomes *inaccessible* if it terminates or if it is *aborted*. An active partition is aborted when its environment task is aborted. In addition, if a partition fails during its elaboration, it becomes inaccessible to other partitions. Other implementation-defined events can also result in a partition becoming inaccessible.
- 8/1 For a prefix D that denotes a library-level declaration, excepting a declaration of or within a declared-pure library unit, the following attribute is defined:
- 9 D'Partition_Id
 - Denotes a value of the type *universal_integer* that identifies the partition in which D was elaborated. If D denotes the declaration of a remote call interface library unit (see E.2.3) the given partition is the one where the body of D was elaborated.

Bounded (Run-Time) Errors

- 10 It is a bounded error for there to be cyclic elaboration dependences between the active partitions of a single distributed program. The possible effects, in each of the partitions involved, are deadlock during elaboration, or the raising of Communication_Error or Program_Error.

Implementation Permissions

- 11 An implementation may allow multiple active or passive partitions to be configured on a single processing node, and multiple passive partitions to be configured on a single storage node. In these cases, the scheduling policies, treatment of priorities, and management of shared resources between these partitions are implementation defined.
- 12 An implementation may allow separate copies of an active partition to be configured on different processing nodes, and to provide appropriate interactions between the copies to present a consistent state of the partition to other active partitions.
- 13 In an implementation, the partitions of a distributed program need not be loaded and elaborated all at the same time; they may be loaded and elaborated one at a time over an extended period of time. An implementation may provide facilities to abort and reload a partition during the execution of a distributed program.
- 14 An implementation may allow the state of some of the partitions of a distributed program to persist while other partitions of the program terminate and are later reinvoked.

NOTES

- 15 3 Library units are grouped into partitions after compile time, but before run time. At compile time, only the relevant library unit properties are identified using categorization pragmas.
- 16 4 The value returned by the Partition_Id attribute can be used as a parameter to implementation-provided subprograms in order to query information about the partition.

E.2 Categorization of Library Units

- 1 Library units can be categorized according to the role they play in a distributed program. Certain restrictions are associated with each category to ensure that the semantics of a distributed program remain close to the semantics for a nondistributed program.

A *categorization pragma* is a library unit pragma (see 10.1.5) that restricts the declarations, child units, or semantic dependences of the library unit to which it applies. A *categorized library unit* is a library unit to which a categorization pragma applies.

The pragmas Shared_Passive, Remote_Types, and Remote_Call_Interface are categorization pragmas. In addition, for the purposes of this Annex, the pragma Pure (see 10.2.1) is considered a categorization pragma.

A library package or generic library package is called a *shared passive* library unit if a Shared_Passive pragma applies to it. A library package or generic library package is called a *remote types* library unit if a Remote_Types pragma applies to it. A library unit is called a *remote call interface* if a Remote_Call_Interface pragma applies to it. A *normal library unit* is one to which no categorization pragma applies.

The various categories of library units and the associated restrictions are described in this clause and its subclauses. The categories are related hierarchically in that the library units of one category can depend semantically only on library units of that category or an earlier one, except that the body of a remote types or remote call interface library unit is unrestricted.

The overall hierarchy (including declared pure) is as follows:

Declared Pure

Can depend only on other declared pure library units;

Shared Passive

Can depend only on other shared passive or declared pure library units;

Remote Types

The declaration of the library unit can depend only on other remote types library units, or one of the above; the body of the library unit is unrestricted;

Remote Call Interface

The declaration of the library unit can depend only on other remote call interfaces, or one of the above; the body of the library unit is unrestricted;

Normal

Unrestricted.

Declared pure and shared passive library units are preelaborated. The declaration of a remote types or remote call interface library unit is required to be preelaborable.

Implementation Requirements

This paragraph was deleted.

13/1

Implementation Permissions

Implementations are allowed to define other categorization pragmas.

14

E.2.1 Shared Passive Library Units

A shared passive library unit is used for managing global data shared between active partitions. The restrictions on shared passive library units prevent the data or tasks of one active partition from being accessible to another active partition through references implicit in objects declared in the shared passive library unit.

Syntax

The form of a pragma Shared_Passive is as follows:

2

3 **pragma Shared_Passive[(library_unit_name)];**

Legality Rules

4 A *shared passive library unit* is a library unit to which a Shared_Passive pragma applies. The following restrictions apply to such a library unit:

- 5 • it shall be preelaborable (see 10.2.1);
- 6 • it shall depend semantically only upon declared pure or shared passive library units;
- 7/1 • it shall not contain a library-level declaration of an access type that designates a class-wide type, task type, or protected type with `entry_declarations`.

8 Notwithstanding the definition of accessibility given in 3.10.2, the declaration of a library unit P1 is not accessible from within the declarative region of a shared passive library unit P2, unless the shared passive library unit P2 depends semantically on P1.

Static Semantics

9 A shared passive library unit is preelaborated.

Post-Compilation Rules

10 A shared passive library unit shall be assigned to at most one partition within a given program.

11 Notwithstanding the rule given in 10.2, a compilation unit in a given partition does not *need* (in the sense of 10.2) the shared passive library units on which it depends semantically to be included in that same partition; they will typically reside in separate passive partitions.

E.2.2 Remote Types Library Units

1 A remote types library unit supports the definition of types intended for use in communication between active partitions.

Syntax

2 The form of a **pragma Remote_Types** is as follows:

3 **pragma Remote_Types[(library_unit_name)];**

Legality Rules

4 A *remote types library unit* is a library unit to which the **pragma Remote_Types** applies. The following restrictions apply to the declaration of such a library unit:

- 5 • it shall be preelaborable;
- 6 • it shall depend semantically only on declared pure, shared passive, or other remote types library units;
- 7 • it shall not contain the declaration of any variable within the visible part of the library unit;
- 8/2 • the full view of each type declared in the visible part of the library unit that has any available stream attributes shall support external streaming (see 13.13.2).

9/1 An access type declared in the visible part of a remote types or remote call interface library unit is called a *remote access type*. Such a type shall be:

- 9.1/1 • an access-to-subprogram type, or
- 9.2/1 • a general access type that designates a class-wide limited private type or a class-wide private type extension all of whose ancestors are either private type extensions or limited private types.

A type that is derived from a remote access type is also a remote access type.	9.3/1
The following restrictions apply to the use of a remote access-to-subprogram type:	10
<ul style="list-style-type: none"> • A value of a remote access-to-subprogram type shall be converted only to or from another (subtype-conformant) remote access-to-subprogram type; • The prefix of an Access attribute reference that yields a value of a remote access-to-subprogram type shall statically denote a (subtype-conformant) remote subprogram. 	11/2
The following restrictions apply to the use of a remote access-to-class-wide type:	12
<ul style="list-style-type: none"> • The primitive subprograms of the corresponding specific limited private type shall only have access parameters if they are controlling formal parameters; each non-controlling formal parameter shall support external streaming (see 13.13.2); • A value of a remote access-to-class-wide type shall be explicitly converted only to another remote access-to-class-wide type; • A value of a remote access-to-class-wide type shall be dereferenced (or implicitly converted to an anonymous access type) only as part of a dispatching call where the value designates a controlling operand of the call (see E.4, “Remote Subprogram Calls”). • The Storage_Pool attribute is not defined for a remote access-to-class-wide type; the expected type for an allocator shall not be a remote access-to-class-wide type. A remote access-to-class-wide type shall not be an actual parameter for a generic formal access type. The Storage_Size attribute of a remote access-to-class-wide type yields 0; it is not allowed in an attribute_definition_clause. 	13
NOTES	14/2
5 A remote types library unit need not be pure, and the types it defines may include levels of indirection implemented by using access types. User-specified Read and Write attributes (see 13.13.2) provide for sending values of such a type between active partitions, with Write marshalling the representation, and Read unmarshalling any levels of indirection.	15
	16/1
	17/2
	18

E.2.3 Remote Call Interface Library Units

A remote call interface library unit can be used as an interface for remote procedure calls (RPCs) (or remote function calls) between active partitions.

Syntax

The form of a pragma Remote_Call_Interface is as follows:

pragma Remote_Call_Interface[(*library_unit_name*)];

The form of a pragma All_Calls_Remote is as follows:

pragma All_Calls_Remote[(*library_unit_name*)];

A pragma All_Calls_Remote is a library unit pragma.

Legality Rules

A *remote call interface (RCI)* is a library unit to which the pragma Remote_Call_Interface applies. A subprogram declared in the visible part of such a library unit, or declared by such a library unit, is called a *remote subprogram*.

The declaration of an RCI library unit shall be preelaborable (see 10.2.1), and shall depend semantically only upon declared pure, shared passive, remote types, or other remote call interface library units.

In addition, the following restrictions apply to an RCI library unit:

- its visible part shall not contain the declaration of a variable;

- 11/1 • its visible part shall not contain the declaration of a limited type;
- 12/1 • its visible part shall not contain a nested generic_declaration;
- 13/1 • it shall not be, nor shall its visible part contain, the declaration of a subprogram to which a pragma Inline applies;
- 14/2 • it shall not be, nor shall its visible part contain, a subprogram (or access-to-subprogram) declaration whose profile has an access parameter or a parameter of a type that does not support external streaming (see 13.13.2);
- 15 • any public child of the library unit shall be a remote call interface library unit.
- 16 If a pragma All_Calls_Remote applies to a library unit, the library unit shall be a remote call interface.

Post-Compilation Rules

- 17 A remote call interface library unit shall be assigned to at most one partition of a given program. A remote call interface library unit whose parent is also an RCI library unit shall be assigned only to the same partition as its parent.
- 18 Notwithstanding the rule given in 10.2, a compilation unit in a given partition that semantically depends on the declaration of an RCI library unit, *needs* (in the sense of 10.2) only the declaration of the RCI library unit, not the body, to be included in that same partition. Therefore, the body of an RCI library unit is included only in the partition to which the RCI library unit is explicitly assigned.

Implementation Requirements

- 19/1 If a pragma All_Calls_Remote applies to a given RCI library unit, then the implementation shall route any call to a subprogram of the RCI unit from outside the declarative region of the unit through the Partition Communication Subsystem (PCS); see E.5. Calls to such subprograms from within the declarative region of the unit are defined to be local and shall not go through the PCS.

Implementation Permissions

- 20 An implementation need not support the Remote_Call_Interface pragma nor the All_Calls_Remote pragma. Explicit message-based communication between active partitions can be supported as an alternative to RPC.

E.3 Consistency of a Distributed System

- 1 This clause defines attributes and rules associated with verifying the consistency of a distributed program.

Static Semantics

- 2/1 For a prefix P that statically denotes a program unit, the following attributes are defined:
- 3 P'Version Yields a value of the predefined type String that identifies the version of the compilation unit that contains the declaration of the program unit.
- 4 P'Body_Version Yields a value of the predefined type String that identifies the version of the compilation unit that contains the body (but not any subunits) of the program unit.
- 5/1 The *version* of a compilation unit changes whenever the compilation unit changes in a semantically significant way. This International Standard does not define the exact meaning of "semantically significant". It is unspecified whether there are other events (such as recompilation) that result in the version of a compilation unit changing.

If P is not a library unit, and P has no completion, then P'Body_Version returns the Body_Version of the innermost program unit enclosing the declaration of P. If P is a library unit, and P has no completion, then P'Body_Version returns a value that is different from Body_Version of any version of P that has a completion.

5.1/1

Bounded (Run-Time) Errors

In a distributed program, a library unit is *consistent* if the same version of its declaration is used throughout. It is a bounded error to elaborate a partition of a distributed program that contains a compilation unit that depends on a different version of the declaration of a shared passive or RCI library unit than that included in the partition to which the shared passive or RCI library unit was assigned. As a result of this error, Program_Error can be raised in one or both partitions during elaboration; in any case, the partitions become inaccessible to one another.

6

E.4 Remote Subprogram Calls

A *remote subprogram call* is a subprogram call that invokes the execution of a subprogram in another partition. The partition that originates the remote subprogram call is the *calling partition*, and the partition that executes the corresponding subprogram body is the *called partition*. Some remote procedure calls are allowed to return prior to the completion of subprogram execution. These are called *asynchronous remote procedure calls*.

1

There are three different ways of performing a remote subprogram call:

2

- As a direct call on a (remote) subprogram explicitly declared in a remote call interface;
- As an indirect call through a value of a remote access-to-subprogram type;
- As a dispatching call with a controlling operand designated by a value of a remote access-to-class-wide type.

3

4

5

The first way of calling corresponds to a *static* binding between the calling and the called partition. The latter two ways correspond to a *dynamic* binding between the calling and the called partition.

6

A remote call interface library unit (see E.2.3) defines the remote subprograms or remote access types used for remote subprogram calls.

7

Legality Rules

In a dispatching call with two or more controlling operands, if one controlling operand is designated by a value of a remote access-to-class-wide type, then all shall be.

8

Dynamic Semantics

For the execution of a remote subprogram call, subprogram parameters (and later the results, if any) are passed using a stream-oriented representation (see 13.13.1) which is suitable for transmission between partitions. This action is called *marshalling*. *Unmarshalling* is the reverse action of reconstructing the parameters or results from the stream-oriented representation. Marshalling is performed initially as part of the remote subprogram call in the calling partition; unmarshalling is done in the called partition. After the remote subprogram completes, marshalling is performed in the called partition, and finally unmarshalling is done in the calling partition.

9

A *calling stub* is the sequence of code that replaces the subprogram body of a remotely called subprogram in the calling partition. A *receiving stub* is the sequence of code (the “wrapper”) that receives a remote subprogram call on the called partition and invokes the appropriate subprogram body.

10

- 11 Remote subprogram calls are executed at most once, that is, if the subprogram call returns normally, then the called subprogram's body was executed exactly once.
- 12 The task executing a remote subprogram call blocks until the subprogram in the called partition returns, unless the call is asynchronous. For an asynchronous remote procedure call, the calling task can become ready before the procedure in the called partition returns.
- 13 If a construct containing a remote call is aborted, the remote subprogram call is *cancelled*. Whether the execution of the remote subprogram is immediately aborted as a result of the cancellation is implementation defined.
- 14 If a remote subprogram call is received by a called partition before the partition has completed its elaboration, the call is kept pending until the called partition completes its elaboration (unless the call is cancelled by the calling partition prior to that).
- 15 If an exception is propagated by a remotely called subprogram, and the call is not an asynchronous call, the corresponding exception is reraised at the point of the remote subprogram call. For an asynchronous call, if the remote procedure call returns prior to the completion of the remotely called subprogram, any exception is lost.
- 16 The exception `Communication_Error` (see E.5) is raised if a remote call cannot be completed due to difficulties in communicating with the called partition.
- 17 All forms of remote subprogram calls are potentially blocking operations (see 9.5.1).
- 18/1 In a remote subprogram call with a formal parameter of a class-wide type, a check is made that the tag of the actual parameter identifies a tagged type declared in a declared-pure or shared passive library unit, or in the visible part of a remote types or remote call interface library unit. `Program_Error` is raised if this check fails. In a remote function call which returns a class-wide type, the same check is made on the function result.
- 19 In a dispatching call with two or more controlling operands that are designated by values of a remote access-to-class-wide type, a check is made (in addition to the normal `Tag_Check` — see 11.5) that all the remote access-to-class-wide values originated from Access `attribute_references` that were evaluated by tasks of the same active partition. `Constraint_Error` is raised if this check fails.

Implementation Requirements

- 20 The implementation of remote subprogram calls shall conform to the PCS interface as defined by the specification of the language-defined package `System.RPC` (see E.5). The calling stub shall use the `Do_RPC` procedure unless the remote procedure call is asynchronous in which case `Do_APC` shall be used. On the receiving side, the corresponding receiving stub shall be invoked by the RPC-receiver.
- 20.1/1 With respect to shared variables in shared passive library units, the execution of the corresponding subprogram body of a synchronous remote procedure call is considered to be part of the execution of the calling task. The execution of the corresponding subprogram body of an asynchronous remote procedure call proceeds in parallel with the calling task and does not signal the next action of the calling task (see 9.10).

NOTES

- 21 6 A given active partition can both make and receive remote subprogram calls. Thus, an active partition can act as both a client and a server.
- 22 7 If a given exception is propagated by a remote subprogram call, but the exception does not exist in the calling partition, the exception can be handled by an `others` choice or be propagated to and handled by a third partition.

E.4.1 Pragma Asynchronous

This subclause introduces the pragma Asynchronous which allows a remote subprogram call to return prior to completion of the execution of the corresponding remote subprogram body.

Syntax

The form of a pragma Asynchronous is as follows:

pragma Asynchronous(*local_name*);

Legality Rules

The *local_name* of a pragma Asynchronous shall denote either:

- One or more remote procedures; the formal parameters of the procedure(s) shall all be of mode **in**;
- The first subtype of a remote access-to-procedure type; the formal parameters of the designated profile of the type shall all be of mode **in**;
- The first subtype of a remote access-to-class-wide type.

Static Semantics

A pragma Asynchronous is a representation pragma. When applied to a type, it specifies the type-related *asynchronous* aspect of the type.

Dynamic Semantics

A remote call is *asynchronous* if it is a call to a procedure, or a call through a value of an access-to-procedure type, to which a pragma Asynchronous applies. In addition, if a pragma Asynchronous applies to a remote access-to-class-wide type, then a dispatching call on a procedure with a controlling operand designated by a value of the type is asynchronous if the formal parameters of the procedure are all of mode **in**.

Implementation Requirements

Asynchronous remote procedure calls shall be implemented such that the corresponding body executes at most once as a result of the call.

E.4.2 Example of Use of a Remote Access-to-Class-Wide Type

Examples

Example of using a remote access-to-class-wide type to achieve dynamic binding across active partitions:

```
package Tapes is
  pragma Pure(Tapes);
  type Tape is abstract tagged limited private;
  -- Primitive dispatching operations where
  -- Tape is controlling operand
  procedure Copy (From, To : access Tape; Num_Recs : in Natural) is
    abstract;
    procedure Rewind (T : access Tape) is abstract;
    -- More operations
  private
    type Tape is ...
  end Tapes;
```

```

3      with Tapes;
4      package Name_Server is
5          pragma Remote_Call_Interface;
6          -- Dynamic binding to remote operations is achieved
7          -- using the access-to-limited-class-wide type Tape_Ptr
8          type Tape_Ptr is access all Tapes.Tape'Class;
9          -- The following statically bound remote operations
10         -- allow for a name-server capability in this example
11         function Find    (Name : String) return Tape_Ptr;
12         procedure Register (Name : in String; T : in Tape_Ptr);
13         procedure Remove   (T : in Tape_Ptr);
14         -- More operations
15     end Name_Server;
16
17     package Tape_Driver is
18         -- Declarations are not shown, they are irrelevant here
19     end Tape_Driver;
20
21     with Tapes, Name_Server;
22     package body Tape_Driver is
23         type New_Tape is new Tapes.Tape with ...
24         procedure Copy
25             (From, To : access New_Tape; Num_Recs: in Natural) is
26         begin
27
28             end Copy;
29             procedure Rewind (T : access New_Tape) is
30             begin
31
32                 end Rewind;
33                 -- Objects remotely accessible through use
34                 -- of Name_Server operations
35                 Tape1, Tape2 : aliased New_Tape;
36
37             begin
38                 Name_Server.Register ("NINE-TRACK", Tape1'Access);
39                 Name_Server.Register ("SEVEN-TRACK", Tape2'Access);
40             end Tape_Driver;
41
42             with Tapes, Name_Server;
43             -- Tape_Driver is not needed and thus not mentioned in the with_clause
44             procedure Tape_Client is
45                 T1, T2 : Name_Server.Tape_Ptr;
46
47             begin
48                 T1 := Name_Server.Find ("NINE-TRACK");
49                 T2 := Name_Server.Find ("SEVEN-TRACK");
50                 Tapes.Rewind (T1);
51                 Tapes.Rewind (T2);
52                 Tapes.Copy (T1, T2, 3);
53             end Tape_Client;

```

7 *Notes on the example:*

8/1 This paragraph was deleted.

- 9 • The package Tapes provides the necessary declarations of the type and its primitive operations.
- 10 • Name_Server is a remote call interface package and is elaborated in a separate active partition to provide the necessary naming services (such as Register and Find) to the entire distributed program through remote subprogram calls.
- 11 • Tape_Driver is a normal package that is elaborated in a partition configured on the processing node that is connected to the tape device(s). The abstract operations are overridden to support the locally declared tape devices (Tape1, Tape2). The package is not visible to its clients, but it exports the tape devices (as remote objects) through the services of the Name_Server. This allows for tape devices to be dynamically added, removed or replaced without requiring the modification of the clients' code.

- The Tape_Client procedure references only declarations in the Tapes and Name_Server packages. Before using a tape for the first time, it needs to query the Name_Server for a system-wide identity for that tape. From then on, it can use that identity to access the tape device. 12
- Values of remote access type Tape_Ptr include the necessary information to complete the remote dispatching operations that result from dereferencing the controlling operands T1 and T2. 13

E.5 Partition Communication Subsystem

The *Partition Communication Subsystem* (PCS) provides facilities for supporting communication between the active partitions of a distributed program. The package System.RPC is a language-defined interface to the PCS. 1/2

Static Semantics

The following language-defined library package exists: 2

```

with Ada.Streams; -- see 13.13.1
package System.RPC is
    type Partition_Id is range 0 .. implementation-defined;
    Communication_Error : exception;
    type Params_Stream_Type (
        Initial_Size : Ada.Streams.Stream_Element_Count) is new
        Ada.Streams.Root_Stream_Type with private;
    procedure Read(
        Stream : in out Params_Stream_Type;
        Item : out Ada.Streams.Stream_Element_Array;
        Last : out Ada.Streams.Stream_Element_Offset);
    procedure Write(
        Stream : in out Params_Stream_Type;
        Item : in Ada.Streams.Stream_Element_Array);
    -- Synchronous call
    procedure Do_RPC(
        Partition : in Partition_Id;
        Params    : access Params_Stream_Type;
        Result    : access Params_Stream_Type);
    -- Asynchronous call
    procedure Do_APC(
        Partition : in Partition_Id;
        Params    : access Params_Stream_Type);
    -- The handler for incoming RPCs
    type RPC_Receiver is access procedure(
        Params    : access Params_Stream_Type;
        Result    : access Params_Stream_Type);
    procedure Establish_RPC_Receiver(
        Partition : in Partition_Id;
        Receiver  : in RPC_Receiver);
private
    ... -- not specified by the language
end System.RPC;

```

A value of the type Partition_Id is used to identify a partition. 14

An object of the type Params_Stream_Type is used for identifying the particular remote subprogram that is being called, as well as marshalling and unmarshalling the parameters or result of a remote subprogram call, as part of sending them between partitions. 15

- 16 The Read and Write procedures override the corresponding abstract operations for the type Params_Stream_Type.

Dynamic Semantics

- 17 The Do_RPC and Do_APC procedures send a message to the active partition identified by the Partition parameter.
- 18 After sending the message, Do_RPC blocks the calling task until a reply message comes back from the called partition or some error is detected by the underlying communication system in which case Communication_Error is raised at the point of the call to Do_RPC.
- 19 Do_APC operates in the same way as Do_RPC except that it is allowed to return immediately after sending the message.
- 20 Upon normal return, the stream designated by the Result parameter of Do_RPC contains the reply message.
- 21 The procedure System.RPC.Establish_RPC_Receiver is called once, immediately after elaborating the library units of an active partition (that is, right after the *elaboration of the partition*) if the partition includes an RCI library unit, but prior to invoking the main subprogram, if any. The Partition parameter is the Partition_Id of the active partition being elaborated. The Receiver parameter designates an implementation-provided procedure called the *RPC-receiver* which will handle all RPCs received by the partition from the PCS. Establish_RPC_Receiver saves a reference to the RPC-receiver; when a message is received at the called partition, the RPC-receiver is called with the Params stream containing the message. When the RPC-receiver returns, the contents of the stream designated by Result is placed in a message and sent back to the calling partition.
- 22 If a call on Do_RPC is aborted, a cancellation message is sent to the called partition, to request that the execution of the remotely called subprogram be aborted.
- 23 The subprograms declared in System.RPC are potentially blocking operations.

Implementation Requirements

- 24 The implementation of the RPC-receiver shall be reentrant, thereby allowing concurrent calls on it from the PCS to service concurrent remote subprogram calls into the partition.
- 24.1/1 An implementation shall not restrict the replacement of the body of System.RPC. An implementation shall not restrict children of System.RPC. The related implementation permissions in the introduction to Annex A do not apply.
- 24.2/1 If the implementation of System.RPC is provided by the user, an implementation shall support remote subprogram calls as specified.

Documentation Requirements

- 25 The implementation of the PCS shall document whether the RPC-receiver is invoked from concurrent tasks. If there is an upper limit on the number of such tasks, this limit shall be documented as well, together with the mechanisms to configure it (if this is supported).

Implementation Permissions

- 26 The PCS is allowed to contain implementation-defined interfaces for explicit message passing, broadcasting, etc. Similarly, it is allowed to provide additional interfaces to query the state of some remote

partition (given its partition ID) or of the PCS itself, to set timeouts and retry parameters, to get more detailed error status, etc. These additional interfaces should be provided in child packages of System.RPC.

A body for the package System.RPC need not be supplied by the implementation. 27

An alternative declaration is allowed for package System.RPC as long as it provides a set of operations 27.1/2 that is substantially equivalent to the specification defined in this clause.

Implementation Advice

Whenever possible, the PCS on the called partition should allow for multiple tasks to call the RPC-receiver with different messages and should allow them to block until the corresponding subprogram body returns. 28

The Write operation on a stream of type Params_Stream_Type should raise Storage_Error if it runs out of space trying to write the Item into the stream. 29

NOTES

8 The package System.RPC is not designed for direct calls by user programs. It is instead designed for use in the implementation of remote subprograms calls, being called by the calling stubs generated for a remote call interface library unit to initiate a remote call, and in turn calling back to an RPC-receiver that dispatches to the receiving stubs generated for the body of a remote call interface, to handle a remote call received from elsewhere. 30

Annex F (normative) Information Systems

This Annex provides a set of facilities relevant to Information Systems programming. These fall into several categories:

- an attribute definition clause specifying Machine_Radix for a decimal subtype; 1
- the package Decimal, which declares a set of constants defining the implementation's capacity 2 for decimal types, and a generic procedure for decimal division; and 3
- the child packages Text_IO.Editing, Wide_Text_IO.Editing, and Wide_Wide_Text_IO.Editing, 4/2 which support formatted and localized output of decimal data, based on "picture String" values.

See also: 3.5.9, "Fixed Point Types"; 3.5.10, "Operations of Fixed Point Types"; 4.6, "Type Conversions"; 5/2 13.3, "Operational and Representation Attributes"; A.10.9, "Input-Output for Real Types"; B.3, "Interfacing with C and C++"; B.4, "Interfacing with COBOL"; Annex G, "Numerics".

The character and string handling packages in Annex A, "Predefined Language Environment" are also 6 relevant for Information Systems.

Implementation Advice

If COBOL (respectively, C) is widely supported in the target environment, implementations supporting the 7 Information Systems Annex should provide the child package Interfaces.COBOL (respectively, Interfaces.C) specified in Annex B and should support a *convention_identifier* of COBOL (respectively, C) in the interfacing pragmas (see Annex B), thus allowing Ada programs to interface with programs written in that language.

F.1 Machine_Radix Attribute Definition Clause

Static Semantics

Machine_Radix may be specified for a decimal first subtype (see 3.5.9) via an *attribute_definition_clause*; 1 the expression of such a clause shall be static, and its value shall be 2 or 10. A value of 2 implies a binary base range; a value of 10 implies a decimal base range.

Implementation Advice

Packed decimal should be used as the internal representation for objects of subtype S when 2 S'Machine_Radix = 10.

Examples

Example of Machine_Radix attribute definition clause:

```
type Money is delta 0.01 digits 15;
for Money'Machine_Radix use 10;
```

F.2 The Package Decimal

Static Semantics

1 The library package `Decimal` has the following declaration:

```

2   package Ada.Decimal is
3     pragma Pure(Decimal);
4     Max_Scale : constant := implementation-defined;
5     Min_Scale : constant := implementation-defined;
6     Min_Delta : constant := 10.0**(-Max_Scale);
7     Max_Delta : constant := 10.0**(-Min_Scale);
8     Max_Decimal_Digits : constant := implementation-defined;
9
10    generic
11      type Dividend_Type is delta <> digits <>;
12      type Divisor_Type is delta <> digits <>;
13      type Quotient_Type is delta <> digits <>;
14      type Remainder_Type is delta <> digits <>;
15      procedure Divide (Dividend : in Dividend_Type;
16                        Divisor : in Divisor_Type;
17                        Quotient : out Quotient_Type;
18                        Remainder : out Remainder_Type);
19    pragma Convention(Intrinsic, Divide);
20
21  end Ada.Decimal;
```

8 `Max_Scale` is the largest N such that $10.0^{**(-N)}$ is allowed as a decimal type's delta. Its type is *universal_integer*.

9 `Min_Scale` is the smallest N such that $10.0^{**(-N)}$ is allowed as a decimal type's delta. Its type is *universal_integer*.

10 `Min_Delta` is the smallest value allowed for *delta* in a `decimal_fixed_point_definition`. Its type is *universal_real*.

11 `Max_Delta` is the largest value allowed for *delta* in a `decimal_fixed_point_definition`. Its type is *universal_real*.

12 `Max_Decimal_Digits` is the largest value allowed for *digits* in a `decimal_fixed_point_definition`. Its type is *universal_integer*.

Static Semantics

13 The effect of `Divide` is as follows. The value of `Quotient` is `Quotient_Type(Dividend/Divisor)`. The value of `Remainder` is `Remainder_Type(Intermediate)`, where `Intermediate` is the difference between `Dividend` and the product of `Divisor` and `Quotient`; this result is computed exactly.

Implementation Requirements

14 `Decimal.Max_Decimal_Digits` shall be at least 18.

15 `Decimal.Max_Scale` shall be at least 18.

16 `Decimal.Min_Scale` shall be at most 0.

NOTES

17 1 The effect of division yielding a quotient with control over rounding versus truncation is obtained by applying either the function attribute `Quotient_Type'Round` or the conversion `Quotient_Type` to the expression `Dividend/Divisor`.

F.3 Edited Output for Decimal Types

The child packages `Text_IO.Editing`, `Wide_Text_IO.Editing`, and `Wide_Wide_Text_IO.Editing` provide localizable formatted text output, known as *edited output*, for decimal types. An edited output string is a function of a numeric value, program-specifiable locale elements, and a format control value. The numeric value is of some decimal type. The locale elements are:

- the currency string;
- the digits group separator character;
- the radix mark character; and
- the fill character that replaces leading zeros of the numeric value.

For `Text_IO.Editing` the edited output and currency strings are of type `String`, and the locale characters are of type `Character`. For `Wide_Text_IO.Editing` their types are `Wide_String` and `Wide_Character`, respectively. For `Wide_Wide_Text_IO.Editing` their types are `Wide_Wide_String` and `Wide_Wide_Character`, respectively.

Each of the locale elements has a default value that can be replaced or explicitly overridden.

A format-control value is of the private type `Picture`; it determines the composition of the edited output string and controls the form and placement of the sign, the position of the locale elements and the decimal digits, the presence or absence of a radix mark, suppression of leading zeros, and insertion of particular character values.

A `Picture` object is composed from a `String` value, known as a *picture String*, that serves as a template for the edited output string, and a Boolean value that controls whether a string of all space characters is produced when the number's value is zero. A picture String comprises a sequence of one- or two-Character symbols, each serving as a placeholder for a character or string at a corresponding position in the edited output string. The picture String symbols fall into several categories based on their effect on the edited output string:

Decimal Digit:	'9'					
Radix Control:	'.'	'V'				
Sign Control:	'+'	'-'	'<'	'>'	"CR"	"DB"
Currency Control:	'\$'	'#'				
Zero Suppression:	'Z'	'*''				
Simple Insertion:	'_'	'B'	'0'	'/'		

The entries are not case-sensitive. Mixed- or lower-case forms for "CR" and "DB", and lower-case forms for 'V', 'Z', and 'B', have the same effect as the upper-case symbols shown.

An occurrence of a '9' Character in the picture String represents a decimal digit position in the edited output string.

A radix control Character in the picture String indicates the position of the radix mark in the edited output string: an actual character position for '.', or an assumed position for 'V'.

A sign control Character in the picture String affects the form of the sign in the edited output string. The '<' and '>' Character values indicate parentheses for negative values. A Character '+', '−', or '<' appears either singly, signifying a fixed-position sign in the edited output, or repeated, signifying a floating-position sign that is preceded by zero or more space characters and that replaces a leading 0.

- 15 A currency control Character in the picture String indicates an occurrence of the currency string in the edited output string. The '\$' Character represents the complete currency string; the '#' Character represents one character of the currency string. A '\$' Character appears either singly, indicating a fixed-position currency string in the edited output, or repeated, indicating a floating-position currency string that occurs in place of a leading 0. A sequence of '#' Character values indicates either a fixed- or floating-position currency string, depending on context.
- 16 A zero suppression Character in the picture String allows a leading zero to be replaced by either the space character (for 'Z') or the fill character (for '*').
- 17 A simple insertion Character in the picture String represents, in general, either itself (if '/' or '0'), the space character (if 'B'), or the digits group separator character (if '_'). In some contexts it is treated as part of a floating sign, floating currency, or zero suppression string.
- 18/2 An example of a picture String is "<###Z_ZZ9.99>". If the currency string is "kr", the separator character is ',', and the radix mark is '.' then the edited output string values for the decimal values 32.10 and –5432.10 are "bbkrbbb32.10b" and "(bkr5,432.10)", respectively, where 'b' indicates the space character.
- 19/2 The generic packages `Text_IO.Decimal_IO`, `Wide_Text_IO.Decimal_IO`, and `Wide_Wide_Text_IO.Decimal_IO` (see A.10.9, “Input-Output for Real Types”) provide text input and non-edited text output for decimal types.

NOTES

- 20/2 2 A picture String is of type `Standard.String`, for all of `Text_IO.Editing`, `Wide_Text_IO.Editing`, and `Wide_Wide_Text_IO.Editing`.

F.3.1 Picture String Formation

- 1 A *well-formed picture String*, or simply *picture String*, is a String value that conforms to the syntactic rules, composition constraints, and character replication conventions specified in this clause.

Dynamic Semantics

- 2/1 This paragraph was deleted.

```
3 picture_string ::=
  fixed_$_picture_string
  | fixed_#_picture_string
  | floating_currency_picture_string
  | non_currency_picture_string
```

fixed_\$_picture_string ::=
 [fixed_LHS_sign] fixed\$_char {direct_insertion} [zero_suppression]
 number [RHS_sign]
 | [fixed_LHS_sign {direct_insertion}] [zero_suppression]
 number fixed\$_char {direct_insertion} [RHS_sign]
 | floating_LHS_sign number fixed\$_char {direct_insertion} [RHS_sign]
 | [fixed_LHS_sign] fixed\$_char {direct_insertion}
 all_zero_suppression_number {direct_insertion} [RHS_sign]
 | [fixed_LHS_sign {direct_insertion}] all_zero_suppression_number {direct_insertion}
 fixed\$_char {direct_insertion} [RHS_sign]
 | all_sign_number {direct_insertion} fixed\$_char {direct_insertion} [RHS_sign]

fixed_#picture_string ::=
 [fixed_LHS_sign] single_#_currency {direct_insertion}
 [zero_suppression] number [RHS_sign]
 | [fixed_LHS_sign] multiple_#_currency {direct_insertion}
 zero_suppression number [RHS_sign]
 | [fixed_LHS_sign {direct_insertion}] [zero_suppression]
 number fixed_#_currency {direct_insertion} [RHS_sign]
 | floating_LHS_sign number fixed_#_currency {direct_insertion} [RHS_sign]
 | [fixed_LHS_sign] single_#_currency {direct_insertion}
 all_zero_suppression_number {direct_insertion} [RHS_sign]
 | [fixed_LHS_sign] multiple_#_currency {direct_insertion}
 all_zero_suppression_number {direct_insertion} [RHS_sign]
 | [fixed_LHS_sign {direct_insertion}] all_zero_suppression_number {direct_insertion}
 fixed_#_currency {direct_insertion} [RHS_sign]
 | all_sign_number {direct_insertion} fixed_#_currency {direct_insertion} [RHS_sign]

floating_currency_picture_string ::=
 [fixed_LHS_sign] {direct_insertion} floating\$_currency number [RHS_sign]
 | [fixed_LHS_sign] {direct_insertion} floating_#_currency number [RHS_sign]
 | [fixed_LHS_sign] {direct_insertion} all_currency_number {direct_insertion} [RHS_sign]

```

7    non_currency_picture_string ::=  

     | [fixed_LHS_sign {direct_insertion}] zero_suppression number [RHS_sign]  

     | [floating_LHS_sign] number [RHS_sign]  

     | [fixed_LHS_sign {direct_insertion}] all_zero_suppression_number {direct_insertion}  

       [RHS_sign]  

     | all_sign_number {direct_insertion}  

     | fixed_LHS_sign direct_insertion {direct_insertion} number [RHS_sign]  

  

8    fixed_LHS_sign ::= LHS_Sign  

9    LHS_Sign ::= + | - | <  

  

10   fixed_$_char ::= $  

  

11   direct_insertion ::= simple_insertion  

12   simple_insertion ::= _ | B | 0 | /  

  

13   zero_suppression ::= Z {Z | context_sensitive_insertion} | fill_string  

14   context_sensitive_insertion ::= simple_insertion  

  

15   fill_string ::= * {* | context_sensitive_insertion}  

  

16   number ::=  

     | fore_digits [radix [aft_digits] {direct_insertion}]  

     | radix aft_digits {direct_insertion}  

17   fore_digits ::= 9 {9 | direct_insertion}  

18   aft_digits ::= {9 | direct_insertion} 9  

19   radix ::= . | V  

  

20   RHS_sign ::= + | - | > | CR | DB  

  

21   floating_LHS_sign ::=  

     LHS_Sign {context_sensitive_insertion} LHS_Sign {LHS_Sign | context_sensitive_insertion}  

  

22   single_##_currency ::= #  

23   multiple_##_currency ::= ## {#}  

  

24   fixed_##_currency ::= single_##_currency | multiple_##_currency  

  

25   floating_$_currency ::=  

     $ {context_sensitive_insertion} $ {$ | context_sensitive_insertion}  

  

26   floating_##_currency ::=  

     # {context_sensitive_insertion} # {# | context_sensitive_insertion}  

  

27   all_sign_number ::= all_sign_fore [radix [all_sign_aft]] [>]

```

all_sign_fore ::=	28
sign_char {context_sensitive_insertion} sign_char {sign_char context_sensitive_insertion}	
all_sign_aft ::= {all_sign_aft_char} sign_char	29
all_sign_aft_char ::= sign_char context_sensitive_insertion	
sign_char ::= + - <	30
all_currency_number ::= all_currency_fore [radix [all_currency_aft]]	31
all_currency_fore ::=	
currency_char {context_sensitive_insertion}	32
currency_char {currency_char context_sensitive_insertion}	
all_currency_aft ::= {all_currency_aft_char} currency_char	33
all_currency_aft_char ::= currency_char context_sensitive_insertion	
currency_char ::= \$ #	34
all_zero_suppression_number ::= all_zero_suppression_fore [radix [all_zero_suppression_aft]]	35
all_zero_suppression_fore ::=	
zero_suppression_char {zero_suppression_char context_sensitive_insertion}	36
all_zero_suppression_aft ::= {all_zero_suppression_aft_char} zero_suppression_char	37
all_zero_suppression_aft_char ::= zero_suppression_char context_sensitive_insertion	
zero_suppression_char ::= Z *	38

The following composition constraints apply to a picture String:

- A floating_LHS_sign does not have occurrences of different LHS_Sign Character values.
- If a picture String has '<' as fixed_LHS_sign, then it has '>' as RHS_sign.
- If a picture String has '<' in a floating_LHS_sign or in an all_sign_number, then it has an occurrence of '>'.
- If a picture String has '+' or '-' as fixed_LHS_sign, in a floating_LHS_sign, or in an all_sign_number, then it has no RHS_sign or '>' character.
- An instance of all_sign_number does not have occurrences of different sign_char Character values.
- An instance of all_currency_number does not have occurrences of different currency_char Character values.
- An instance of all_zero_suppression_number does not have occurrences of different zero_suppression_char Character values, except for possible case differences between 'Z' and 'z'.

A *replicable Character* is a Character that, by the above rules, can occur in two consecutive positions in a picture String.

A *Character replication* is a String

char & ' (' & spaces & count_string & ') '

where *char* is a replicable Character, *spaces* is a String (possibly empty) comprising only space Character values, and *count_string* is a String of one or more decimal digit Character values. A Character replication

in a picture String has the same effect as (and is said to be *equivalent to*) a String comprising n consecutive occurrences of *char*, where $n = \text{Integer}'\text{Value}(\text{count_string})$.

- 51 An *expanded picture String* is a picture String containing no Character replications.

NOTES

- 52 3 Although a sign to the left of the number can float, a sign to the right of the number is in a fixed position.

F.3.2 Edited Output Generation

Dynamic Semantics

- 1 The contents of an edited output string are based on:
 - A value, Item, of some decimal type Num,
 - An expanded picture String Pic_String,
 - A Boolean value, Blank_When_Zero,
 - A Currency string,
 - A Fill character,
 - A Separator character, and
 - A Radix_Mark character.
- 9 The combination of a True value for Blank_When_Zero and a '*' character in Pic_String is inconsistent; no edited output string is defined.
- 10 A layout error is identified in the rules below if leading non-zero digits of Item, character values of the Currency string, or a negative sign would be truncated; in such cases no edited output string is defined.
- 11 The edited output string has lower bound 1 and upper bound N where $N = \text{Pic_String}'\text{Length} + \text{Currency_Length_Adjustment} - \text{Radix_Adjustment}$, and
 - Currency_Length_Adjustment = Currency'Length - 1 if there is some occurrence of '\$' in Pic_String, and 0 otherwise.
 - Radix_Adjustment = 1 if there is an occurrence of 'V' or 'v' in Pic_Str, and 0 otherwise.
- 14 Let the magnitude of Item be expressed as a base-10 number $I_p \cdots I_1.F_1 \cdots F_q$, called the *displayed magnitude* of Item, where:
 - $q = \text{Min}(\text{Max}(\text{Num}'\text{Scale}, 0), n)$ where n is 0 if Pic_String has no radix and is otherwise the number of digit positions following radix in Pic_String, where a digit position corresponds to an occurrence of '9', a zero_suppression_char (for an all_zero_suppression_number), a currency_char (for an all_currency_number), or a sign_char (for an all_sign_number).
 - $I_p = 0$ if $p > 0$.
- 17 If $n < \text{Num}'\text{Scale}$, then the above number is the result of rounding (away from 0 if exactly midway between values).
- 18 If Blank_When_Zero = True and the displayed magnitude of Item is zero, then the edited output string comprises all space character values. Otherwise, the picture String is treated as a sequence of instances of syntactic categories based on the rules in F.3.1, and the edited output string is the concatenation of string values derived from these categories according to the following mapping rules.

Table F-1 shows the mapping from a sign control symbol to a corresponding character or string in the edited output. In the columns showing the edited output, a lower-case 'b' represents the space character. If there is no sign control symbol but the value of Item is negative, a layout error occurs and no edited output string is produced.

Table F-1: Edited Output for Sign Control Symbols		
Sign Control Symbol	Edited Output for Non-Negative Number	Edited Output for Negative Number
'+'	'+'	'_'
'_'	'b'	'_'
'<'	'b'	('
'>'	'b'	')'
"CR"	"bb"	"CR"
"DB"	"bb"	"DB"

An instance of `fixed_LHS_sign` maps to a character as shown in Table F-1.

An instance of `fixed$_char` maps to `Currency`.

An instance of `direct_insertion` maps to `Separator` if `direct_insertion = '_'`, and to the `direct_insertion Character` otherwise.

An instance of `number` maps to a string `integer_part & radix_part & fraction_part` where:

- The string for `integer_part` is obtained as follows:
 1. Occurrences of '9' in `fore_digits` of `number` are replaced from right to left with the decimal digit character values for I_1, \dots, I_p , respectively.
 2. Each occurrence of '9' in `fore_digits` to the left of the leftmost '9' replaced according to rule 1 is replaced with '0'.
 3. If p exceeds the number of occurrences of '9' in `fore_digits` of `number`, then the excess leftmost digits are eligible for use in the mapping of an instance of `zero_suppression`, `floating_LHS_sign`, `floating$_currency`, or `floating$_#_currency` to the left of `number`; if there is no such instance, then a layout error occurs and no edited output string is produced.
- The `radix_part` is:
 - "" if `number` does not include a `radix`, if `radix = 'V'`, or if `radix = 'v'`
 - `Radix_Mark` if `number` includes '.' as `radix`
- The string for `fraction_part` is obtained as follows:
 1. Occurrences of '9' in `aft_digits` of `number` are replaced from left to right with the decimal digit character values for F_1, \dots, F_q .
 2. Each occurrence of '9' in `aft_digits` to the right of the rightmost '9' replaced according to rule 1 is replaced by '0'.

An instance of `zero_suppression` maps to the string obtained as follows:

- 35 1. The rightmost 'Z', 'z', or '*' Character values are replaced with the excess digits (if any) from the *integer_part* of the mapping of the number to the right of the *zero_suppression* instance,
- 36 2. A *context_sensitive_insertion* Character is replaced as though it were a *direct_insertion* Character, if it occurs to the right of some 'Z', 'z', or '*' in *zero_suppression* that has been mapped to an excess digit,
- 37 3. Each Character to the left of the leftmost Character replaced according to rule 1 above is replaced by:
- 38 • the space character if the zero suppression Character is 'Z' or 'z', or
- 39 • the Fill character if the zero suppression Character is '*'.
- 40 4. A layout error occurs if some excess digits remain after all 'Z', 'z', and '*' Character values in *zero_suppression* have been replaced via rule 1; no edited output string is produced.

41 An instance of *RHS_sign* maps to a character or string as shown in Table F-1.

42 An instance of *floating_LHS_sign* maps to the string obtained as follows.

- 43 1. Up to all but one of the rightmost *LHS_Sign* Character values are replaced by the excess digits (if any) from the *integer_part* of the mapping of the number to the right of the *floating_LHS_sign* instance.
- 44 2. The next Character to the left is replaced with the character given by the entry in Table F-1 corresponding to the *LHS_Sign* Character.
- 45 3. A *context_sensitive_insertion* Character is replaced as though it were a *direct_insertion* Character, if it occurs to the right of the leftmost *LHS_Sign* character replaced according to rule 1.
- 46 4. Any other Character is replaced by the space character..
- 47 5. A layout error occurs if some excess digits remain after replacement via rule 1; no edited output string is produced.

48 An instance of *fixed_#_currency* maps to the Currency string with n space character values concatenated on the left (if the instance does not follow a *radix*) or on the right (if the instance does follow a *radix*), where n is the difference between the length of the *fixed_#_currency* instance and *Currency'Length*. A layout error occurs if *Currency'Length* exceeds the length of the *fixed_#_currency* instance; no edited output string is produced.

49 An instance of *floating_\$_currency* maps to the string obtained as follows:

- 50 1. Up to all but one of the rightmost '\$' Character values are replaced with the excess digits (if any) from the *integer_part* of the mapping of the number to the right of the *floating_\$_currency* instance.
- 51 2. The next Character to the left is replaced by the Currency string.
- 52 3. A *context_sensitive_insertion* Character is replaced as though it were a *direct_insertion* Character, if it occurs to the right of the leftmost '\$' Character replaced via rule 1.
- 53 4. Each other Character is replaced by the space character.
- 54 5. A layout error occurs if some excess digits remain after replacement by rule 1; no edited output string is produced.

55 An instance of *floating_#_currency* maps to the string obtained as follows:

1. Up to all but one of the rightmost '#' Character values are replaced with the excess digits (if any) from the *integer_part* of the mapping of the number to the right of the *floating_#_currency* instance. 56
2. The substring whose last Character occurs at the position immediately preceding the leftmost Character replaced via rule 1, and whose length is *Currency'Length*, is replaced by the *Currency* string. 57
3. A *context_sensitive_insertion* Character is replaced as though it were a *direct_insertion* Character, if it occurs to the right of the leftmost '#' replaced via rule 1. 58
4. Any other Character is replaced by the space character. 59
5. A layout error occurs if some excess digits remain after replacement rule 1, or if there is no substring with the required length for replacement rule 2; no edited output string is produced. 60

An instance of *all_zero_suppression_number* maps to: 61

- a string of all spaces if the displayed magnitude of Item is zero, the *zero_suppression_char* is 'Z' or 'z', and the instance of *all_zero_suppression_number* does not have a *radix* at its last character position; 62
- a string containing the *Fill* character in each position except for the character (if any) corresponding to *radix*, if *zero_suppression_char* = '*' and the displayed magnitude of Item is zero; 63
- otherwise, the same result as if each *zero_suppression_char* in *all_zero_suppression_aft* were '9', interpreting the instance of *all_zero_suppression_number* as either *zero_suppression_number* (if a *radix* and *all_zero_suppression_aft* are present), or as *zero_suppression* otherwise. 64

An instance of *all_sign_number* maps to: 65

- a string of all spaces if the displayed magnitude of Item is zero and the instance of *all_sign_number* does not have a *radix* at its last character position; 66
- otherwise, the same result as if each *sign_char* in *all_sign_number_aft* were '9', interpreting the instance of *all_sign_number* as either *floating_LHS_sign_number* (if a *radix* and *all_sign_number_aft* are present), or as *floating_LHS_sign* otherwise. 67

An instance of *all_currency_number* maps to: 68

- a string of all spaces if the displayed magnitude of Item is zero and the instance of *all_currency_number* does not have a *radix* at its last character position; 69
- otherwise, the same result as if each *currency_char* in *all_currency_number_aft* were '9', interpreting the instance of *all_currency_number* as *floating_\$_currency_number* or *floating_#_currency_number* (if a *radix* and *all_currency_number_aft* are present), or as *floating_\$_currency* or *floating_#_currency* otherwise. 70

Examples

In the result string values shown below, 'b' represents the space character. 71

Item:	Picture and Result Strings:
123456.78	Picture: "-###** *** **9.99" "bbbb\$***123,456.78" "bbFF***123.456,78" (currency = "FF", separator = '.', radix mark = ',')

```

74/1      123456.78      Picture:  "-$**_***_**9.99"
          Result:   "b$***123,456.78"
                           "bFF***123.456,78" (currency = "FF",
                                         separator = ',',
                                         radix mark = ',')
75        0.0       Picture:  "-$$$$$.$$"
          Result:   "bbbbbbbbbb"
76        0.20      Picture:  "-$$$$$.$$"
          Result:   "bbbbbb$ .20"
77      -1234.565     Picture:  "<<<_<<<.=<###>""
          Result:   "bb(1,234.57DMb)" (currency = "DM")
78      12345.67      Picture:  "###_###_##9.99"
          Result:   "bbCHF12,345.67" (currency = "CHF")

```

F.3.3 The Package Text_IO.Editing

- 1 The package Text_IO.Editing provides a private type Picture with associated operations, and a generic package Decimal_Output. An object of type Picture is composed from a well-formed picture String (see F.3.1) and a Boolean item indicating whether a zero numeric value will result in an edited output string of all space characters. The package Decimal_Output contains edited output subprograms implementing the effects defined in F.3.2.

Static Semantics

- 2 The library package Text_IO.Editing has the following declaration:

```

3  package Ada.Text_IO.Editing is
4    type Picture is private;
5    function Valid (Pic_String      : in String;
6                     Blank_When_Zero : in Boolean := False) return Boolean;
7    function To_Picture (Pic_String      : in String;
8                     Blank_When_Zero : in Boolean := False)
9                     return Picture;
10   function Pic_String      (Pic : in Picture) return String;
11   function Blank_When_Zero (Pic : in Picture) return Boolean;
12   Max_Picture_Length   : constant := implementation_defined;
13   Picture_Error         : exception;
14   Default_Currency      : constant String      := "$";
15   Default_Fill           : constant Character := '*';
16   Default_Separator      : constant Character := ',';
17   Default_Radix_Mark    : constant Character := '.';
18
19  generic
20    type Num is delta <> digits <>;
21    Default_Currency      : in String      := Text_IO.Editing.Default_Currency;
22    Default_Fill           : in Character := Text_IO.Editing.Default_Fill;
23    Default_Separator      : in Character := Text_IO.Editing.Default_Separator;
24    Default_Radix_Mark    : in Character := Text_IO.Editing.Default_Radix_Mark;
25
26  package Decimal_Output is
27    function Length (Pic      : in Picture;
28                      Currency : in String := Default_Currency)
29                      return Natural;
30    function Valid (Item     : in Num;
31                     Pic      : in Picture;
32                     Currency : in String := Default_Currency)
33                     return Boolean;

```

```

function Image (Item      : in Num;
                Pic       : in Picture;
                Currency  : in String   := Default_Currency;
                Fill      : in Character := Default_Fill;
                Separator : in Character := Default_Separator;
                Radix_Mark : in Character := Default_Radix_Mark)
        return String;                                         13

procedure Put (File      : in File_Type;
               Item      : in Num;
               Pic       : in Picture;
               Currency  : in String   := Default_Currency;
               Fill      : in Character := Default_Fill;
               Separator : in Character := Default_Separator;
               Radix_Mark : in Character := Default_Radix_Mark);    14

procedure Put (Item      : in Num;
                Pic       : in Picture;
                Currency  : in String   := Default_Currency;
                Fill      : in Character := Default_Fill;
                Separator : in Character := Default_Separator;
                Radix_Mark : in Character := Default_Radix_Mark);    15

procedure Put (To       : out String;
               Item     : in Num;
               Pic      : in Picture;
               Currency : in String   := Default_Currency;
               Fill     : in Character := Default_Fill;
               Separator: in Character := Default_Separator;
               Radix_Mark : in Character := Default_Radix_Mark);    16

end Decimal_Output;
private
  ... -- not specified by the language
end Ada.Text_IO.Editing;                                     17

```

The exception Constraint_Error is raised if the Image function or any of the Put procedures is invoked with a null string for Currency. 17

```

function Valid (Pic_String  : in String;
                 Blank_When_Zero : in Boolean := False) return Boolean;    18

Valid returns True if Pic_String is a well-formed picture String (see F.3.1) the length of whose expansion does not exceed Max_Picture_Length, and if either Blank_When_Zero is False or Pic_String contains no '*'. 19

function To_Picture (Pic_String  : in String;
                      Blank_When_Zero : in Boolean := False)
    return Picture;                                         20

To_Picture returns a result Picture such that the application of the function Pic_String to this result yields an expanded picture String equivalent to Pic_String, and such that Blank_When_Zero applied to the result Picture is the same value as the parameter Blank_When_Zero. Picture_Error is raised if not Valid(Pic_String, Blank_When_Zero). 21

```

```

function Pic_String (Pic : in Picture) return String;          22

function Blank_When_Zero (Pic : in Picture) return Boolean;

If Pic is To_Picture(String_Item, Boolean_Item) for some String_Item and Boolean_Item, then: 23
  • Pic_String(Pic) returns an expanded picture String equivalent to String_Item and with any lower-case letter replaced with its corresponding upper-case form, and 24
  • Blank_When_Zero(Pic) returns Boolean_Item. 25

```

If Pic_1 and Pic_2 are objects of type Picture, then "="(Pic_1, Pic_2) is True when 26

```

27      • Pic_String(Pic_1) = Pic_String(Pic_2), and
28      • Blank_When_Zero(Pic_1) = Blank_When_Zero(Pic_2).

29      function Length (Pic      : in Picture;
30                      Currency : in String := Default_Currency)
31                      return Natural;
32
33      Length returns Pic_String(Pic)'Length + Currency_Length_Adjustment - Radix_Adjustment
34      where
35          • Currency_Length_Adjustment =
36              • Currency'Length - 1 if there is some occurrence of '$' in Pic_String(Pic), and
37              • 0 otherwise.
38
39          • Radix_Adjustment =
40              • 1 if there is an occurrence of 'V' or 'v' in Pic_Str(Pic), and
41              • 0 otherwise.

42      function Valid (Item      : in Num;
43                      Pic       : in Picture;
44                      Currency : in String := Default_Currency)
45                      return Boolean;
46
47      Valid returns True if Image(Item, Pic, Currency) does not raise Layout_Error, and returns False
48      otherwise.

49      function Image (Item      : in Num;
50                      Pic       : in Picture;
51                      Currency : in String    := Default_Currency;
52                      Fill     : in Character := Default_Fill;
53                      Separator : in Character := Default_Separator;
54                      Radix_Mark : in Character := Default_Radix_Mark)
55                      return String;
56
57      Image returns the edited output String as defined in F.3.2 for Item, Pic_String(Pic),
58      Blank_When_Zero(Pic), Currency, Fill, Separator, and Radix_Mark. If these rules identify a
59      layout error, then Image raises the exception Layout_Error.

60      procedure Put (File      : in File_Type;
61                      Item      : in Num;
62                      Pic       : in Picture;
63                      Currency : in String    := Default_Currency;
64                      Fill     : in Character := Default_Fill;
65                      Separator : in Character := Default_Separator;
66                      Radix_Mark : in Character := Default_Radix_Mark);

67      procedure Put (Item      : in Num;
68                      Pic       : in Picture;
69                      Currency : in String    := Default_Currency;
70                      Fill     : in Character := Default_Fill;
71                      Separator : in Character := Default_Separator;
72                      Radix_Mark : in Character := Default_Radix_Mark);

```

42 Each of these Put procedures outputs Image(Item, Pic, Currency, Fill, Separator, Radix_Mark) consistent with the conventions for Put for other real types in case of bounded line length (see A.10.6, “Get and Put Procedures”).

```
procedure Put (To      : out String;
              Item    : in Num;
              Pic     : in Picture;
              Currency : in String := Default_Currency;
              Fill    : in Character := Default_Fill;
              Separator : in Character := Default_Separator;
              Radix_Mark : in Character := Default_Radix_Mark);
```

Put copies Image(Item, Pic, Currency, Fill, Separator, Radix_Mark) to the given string, right justified. Otherwise unassigned Character values in To are assigned the space character. If To'Length is less than the length of the string resulting from Image, then Layout_Error is raised.

Implementation Requirements

Max_Picture_Length shall be at least 30. The implementation shall support currency strings of length up to at least 10, both for Default_Currency in an instantiation of Decimal_Output, and for Currency in an invocation of Image or any of the Put procedures.

NOTES

4 The rules for edited output are based on COBOL (ANSI X3.23:1985, endorsed by ISO as ISO 1989-1985), with the following differences:

- The COBOL provisions for picture string localization and for 'P' format are absent from Ada.
- The following Ada facilities are not in COBOL:
 - currency symbol placement after the number,
 - localization of edited output string for multi-character currency string values, including support for both length-preserving and length-expanding currency symbols in picture strings
 - localization of the radix mark, digits separator, and fill character, and
 - parenthesization of negative values.

The value of 30 for Max_Picture_Length is the same limit as in COBOL.

F.3.4 The Package Wide_Text_IO.Editing

Static Semantics

The child package Wide_Text_IO.Editing has the same contents as Text_IO.Editing, except that:

- each occurrence of Character is replaced by Wide_Character,
- each occurrence of Text_IO is replaced by Wide_Text_IO,
- the subtype of Default_Currency is Wide_String rather than String, and
- each occurrence of String in the generic package Decimal_Output is replaced by Wide_String.

NOTES

5 Each of the functions Wide_Text_IO.Editing.Valid, To_Picture, and Pic_String has String (versus Wide_String) as its parameter or result subtype, since a picture String is not localizable.

F.3.5 The Package Wide_Wide_Text_IO.Editing

Static Semantics

The child package Wide_Wide_Text_IO.Editing has the same contents as Text_IO.Editing, except that:

- each occurrence of Character is replaced by Wide_Wide_Character,
- each occurrence of Text_IO is replaced by Wide_Wide_Text_IO,
- the subtype of Default_Currency is Wide_Wide_String rather than String, and

- 5/2 • each occurrence of String in the generic package Decimal_Output is replaced by Wide_Wide_String.
- NOTES
6/2 6 Each of the functions Wide_Wide_Text_IO.Editing.Valid, To_Picture, and Pic_String has String (versus Wide_Wide_String) as its parameter or result subtype, since a picture String is not localizable.

Annex G (normative) Numerics

The Numerics Annex specifies

- features for complex arithmetic, including complex I/O; 1
- a mode (“strict mode”), in which the predefined arithmetic operations of floating point and fixed point types and the functions and operations of various predefined packages have to provide guaranteed accuracy or conform to other numeric performance requirements, which the Numerics Annex also specifies; 2
- a mode (“relaxed mode”), in which no accuracy or other numeric performance requirements need be satisfied, as for implementations not conforming to the Numerics Annex; 3
- models of floating point and fixed point arithmetic on which the accuracy requirements of strict mode are based; 4
- the definitions of the model-oriented attributes of floating point types that apply in the strict mode; and 5/2
- features for the manipulation of real and complex vectors and matrices. 6.1/2

Implementation Advice

If Fortran (respectively, C) is widely supported in the target environment, implementations supporting the Numerics Annex should provide the child package Interfaces.Fortran (respectively, Interfaces.C) specified in Annex B and should support a *convention identifier* of Fortran (respectively, C) in the interfacing pragmas (see Annex B), thus allowing Ada programs to interface with programs written in that language.

G.1 Complex Arithmetic

Types and arithmetic operations for complex arithmetic are provided in Generic_Complex_Types, which is defined in G.1.1. Implementation-defined approximations to the complex analogs of the mathematical functions known as the “elementary functions” are provided by the subprograms in Generic_Complex_Elementary_Functions, which is defined in G.1.2. Both of these library units are generic children of the predefined package Numerics (see A.5). Nongeneric equivalents of these generic packages for each of the predefined floating point types are also provided as children of Numerics.

G.1.1 Complex Types

Static Semantics

The generic library package Numerics.Generic_Complex_Types has the following declaration:

```
generic
  type Real is digits <>;
package Ada.Numerics.Generic_Complex_Types is
  pragma Pure(Generic_Complex_Types);
  type Complex is
    record
      Re, Im : Real'Base;
    end record;
```

```

4/2      type Imaginary is private;
5       pragma Preelaborable_Initialization(Imaginary);
6       i : constant Imaginary;
7       j : constant Imaginary;
8
9       function Re (X : Complex)    return Real'Base;
10      function Im (X : Complex)    return Real'Base;
11      function Im (X : Imaginary) return Real'Base;
12
13      procedure Set_Re (X : in out Complex;
14                          Re : in     Real'Base);
15      procedure Set_Im (X : in out Complex;
16                          Im : in     Real'Base);
17      procedure Set_Im (X :        out Imaginary;
18                          Im : in     Real'Base);
19
20      function Compose_From_Cartesian (Re, Im : Real'Base) return Complex;
21      function Compose_From_Cartesian (Re      : Real'Base) return Complex;
22      function Compose_From_Cartesian (Im      : Imaginary) return Complex;
23
24      function Modulus (X      : Complex) return Real'Base;
25      function "abs"   (Right : Complex) return Real'Base renames Modulus;
26
27      function Argument (X      : Complex) return Real'Base;
28      function Argument (X      : Complex;
29                           Cycle : Real'Base) return Real'Base;
30
31      function Compose_From_Polar (Modulus, Argument      : Real'Base)
32                      return Complex;
33      function Compose_From_Polar (Modulus, Argument, Cycle : Real'Base)
34                      return Complex;
35
36      function "+"      (Right : Complex) return Complex;
37      function "-"      (Right : Complex) return Complex;
38      function Conjugate (X      : Complex) return Complex;
39
40      function "+"      (Left, Right : Complex) return Complex;
41      function "-"      (Left, Right : Complex) return Complex;
42      function "*"      (Left, Right : Complex) return Complex;
43      function "/"      (Left, Right : Complex) return Complex;
44
45      function "***"   (Left : Complex; Right : Integer) return Complex;
46
47      function "+"      (Right : Imaginary) return Imaginary;
48      function "-"      (Right : Imaginary) return Imaginary;
49      function Conjugate (X      : Imaginary) return Imaginary renames "-";
50      function "abs"   (Right : Imaginary) return Real'Base;
51
52      function "+"      (Left, Right : Imaginary) return Imaginary;
53      function "-"      (Left, Right : Imaginary) return Imaginary;
54      function "*"      (Left, Right : Imaginary) return Real'Base;
55      function "/"      (Left, Right : Imaginary) return Real'Base;
56
57      function "***"   (Left : Imaginary; Right : Integer) return Complex;
58
59      function "<"     (Left, Right : Imaginary) return Boolean;
60      function "<="    (Left, Right : Imaginary) return Boolean;
61      function ">"     (Left, Right : Imaginary) return Boolean;
62      function ">="    (Left, Right : Imaginary) return Boolean;
63
64      function "+"      (Left : Complex;   Right : Real'Base) return Complex;
65      function "+"      (Left : Real'Base; Right : Complex)  return Complex;
66      function "-"      (Left : Complex;   Right : Real'Base) return Complex;
67      function "-"      (Left : Real'Base; Right : Complex)  return Complex;
68      function "*"      (Left : Complex;   Right : Real'Base) return Complex;
69      function "*"      (Left : Real'Base; Right : Complex)  return Complex;
70      function "/"      (Left : Complex;   Right : Real'Base) return Complex;
71      function "/"      (Left : Real'Base; Right : Complex)  return Complex;

```

```

function "+" (Left : Complex; Right : Imaginary) return Complex; 20
function "+" (Left : Imaginary; Right : Complex) return Complex;
function "-" (Left : Complex; Right : Imaginary) return Complex;
function "-" (Left : Imaginary; Right : Complex) return Complex;
function "*" (Left : Complex; Right : Imaginary) return Complex;
function "*" (Left : Imaginary; Right : Complex) return Complex;
function "/" (Left : Complex; Right : Imaginary) return Complex;
function "/" (Left : Imaginary; Right : Complex) return Complex;

function "+" (Left : Imaginary; Right : Real'Base) return Complex; 21
function "+" (Left : Real'Base; Right : Imaginary) return Complex;
function "-" (Left : Imaginary; Right : Real'Base) return Complex;
function "-" (Left : Real'Base; Right : Imaginary) return Complex;
function "*" (Left : Imaginary; Right : Real'Base) return Imaginary;
function "*" (Left : Real'Base; Right : Imaginary) return Imaginary;
function "/" (Left : Imaginary; Right : Real'Base) return Imaginary;
function "/" (Left : Real'Base; Right : Imaginary) return Imaginary;

private 22
type Imaginary is new Real'Base;
i : constant Imaginary := 1.0;
j : constant Imaginary := 1.0;
end Ada.Numerics.Generic_Complex_Types; 23

```

24

The library package Numerics.Complex_Types is declared pure and defines the same types, constants, and subprograms as Numerics.Generic_Complex_Types, except that the predefined type Float is systematically substituted for Real'Base throughout. Nongeneric equivalents of Numerics.Generic_Complex_Types for each of the other predefined floating point types are defined similarly, with the names Numerics.Short_Complex_Types, Numerics.Long_Complex_Types, etc.

25/1

Complex is a visible type with Cartesian components.

26/2

Imaginary is a private type; its full type is derived from Real'Base.

27

The arithmetic operations and the Re, Im, Modulus, Argument, and Conjugate functions have their usual mathematical meanings. When applied to a parameter of pure-imaginary type, the “imaginary-part” function Im yields the value of its parameter, as the corresponding real value. The remaining subprograms have the following meanings:

28

- The Set_Re and Set_Im procedures replace the designated component of a complex parameter with the given real value; applied to a parameter of pure-imaginary type, the Set_Im procedure replaces the value of that parameter with the imaginary value corresponding to the given real value.
 - The Compose_From_Cartesian function constructs a complex value from the given real and imaginary components. If only one component is given, the other component is implicitly zero.
 - The Compose_From_Polar function constructs a complex value from the given modulus (radius) and argument (angle). When the value of the parameter Modulus is positive (resp., negative), the result is the complex value represented by the point in the complex plane lying at a distance from the origin given by the absolute value of Modulus and forming an angle measured counterclockwise from the positive (resp., negative) real axis given by the value of the parameter Argument.
- 29
- 30
- 31

When the Cycle parameter is specified, the result of the Argument function and the parameter Argument of the Compose_From_Polar function are measured in units such that a full cycle of revolution has the given value; otherwise, they are measured in radians.

32

The computed results of the mathematically multivalued functions are rendered single-valued by the following conventions, which are meant to imply the principal branch:

33

- 34 • The result of the Modulus function is nonnegative.
- 35 • The result of the Argument function is in the quadrant containing the point in the complex plane represented by the parameter X. This may be any quadrant (I through IV); thus, the range of the Argument function is approximately $-\pi$ to π ($-{\text{Cycle}/2.0}$ to ${\text{Cycle}/2.0}$, if the parameter Cycle is specified). When the point represented by the parameter X lies on the negative real axis, the result approximates
- 36 • π (resp., $-\pi$) when the sign of the imaginary component of X is positive (resp., negative), if Real'Signed_Zeros is True;
- 37 • π , if Real'Signed_Zeros is False.
- 38 • Because a result lying on or near one of the axes may not be exactly representable, the approximation inherent in computing the result may place it in an adjacent quadrant, close to but on the wrong side of the axis.

Dynamic Semantics

- 39 The exception Numerics.Argument_Error is raised by the Argument and Compose_From_Polar functions with specified cycle, signaling a parameter value outside the domain of the corresponding mathematical function, when the value of the parameter Cycle is zero or negative.
- 40 The exception Constraint_Error is raised by the division operator when the value of the right operand is zero, and by the exponentiation operator when the value of the left operand is zero and the value of the exponent is negative, provided that Real'Machine_Overflows is True; when Real'Machine_Overflows is False, the result is unspecified. Constraint_Error can also be raised when a finite result overflows (see G.2.6).

Implementation Requirements

- 41 In the implementation of Numerics.Generic_Complex_Types, the range of intermediate values allowed during the calculation of a final result shall not be affected by any range constraint of the subtype Real.
- 42 In the following cases, evaluation of a complex arithmetic operation shall yield the *prescribed result*, provided that the preceding rules do not call for an exception to be raised:
- 43 • The results of the Re, Im, and Compose_From_Cartesian functions are exact.
 - 44 • The real (resp., imaginary) component of the result of a binary addition operator that yields a result of complex type is exact when either of its operands is of pure-imaginary (resp., real) type.
 - 45 • The real (resp., imaginary) component of the result of a binary subtraction operator that yields a result of complex type is exact when its right operand is of pure-imaginary (resp., real) type.
 - 46 • The real component of the result of the Conjugate function for the complex type is exact.
 - 47 • When the point in the complex plane represented by the parameter X lies on the nonnegative real axis, the Argument function yields a result of zero.
 - 48 • When the value of the parameter Modulus is zero, the Compose_From_Polar function yields a result of zero.
 - 49 • When the value of the parameter Argument is equal to a multiple of the quarter cycle, the result of the Compose_From_Polar function with specified cycle lies on one of the axes. In this case, one of its components is zero, and the other has the magnitude of the parameter Modulus.
 - 50 • Exponentiation by a zero exponent yields the value one. Exponentiation by a unit exponent yields the value of the left operand. Exponentiation of the value one yields the value one. Exponentiation of the value zero yields the value zero, provided that the exponent is nonzero.

When the left operand is of pure-imaginary type, one component of the result of the exponentiation operator is zero.

When the result, or a result component, of any operator of Numerics.Generic_Complex_Types has a mathematical definition in terms of a single arithmetic or relational operation, that result or result component exhibits the accuracy of the corresponding operation of the type Real. 51

Other accuracy requirements for the Modulus, Argument, and Compose_From_Polar functions, and accuracy requirements for the multiplication of a pair of complex operands or for division by a complex operand, all of which apply only in the strict mode, are given in G.2.6. 52

The sign of a zero result or zero result component yielded by a complex arithmetic operation or function is implementation defined when Real'Signed_Zeros is True. 53

Implementation Permissions

The nongeneric equivalent packages may, but need not, be actual instantiations of the generic package for the appropriate predefined type. 54

Implementations may obtain the result of exponentiation of a complex or pure-imaginary operand by repeated complex multiplication, with arbitrary association of the factors and with a possible final complex reciprocation (when the exponent is negative). Implementations are also permitted to obtain the result of exponentiation of a complex operand, but not of a pure-imaginary operand, by converting the left operand to a polar representation; exponentiating the modulus by the given exponent; multiplying the argument by the given exponent; and reconverting to a Cartesian representation. Because of this implementation freedom, no accuracy requirement is imposed on complex exponentiation (except for the prescribed results given above, which apply regardless of the implementation method chosen). 55/2

Implementation Advice

Because the usual mathematical meaning of multiplication of a complex operand and a real operand is that of the scaling of both components of the former by the latter, an implementation should not perform this operation by first promoting the real operand to complex type and then performing a full complex multiplication. In systems that, in the future, support an Ada binding to IEC 559:1989, the latter technique will not generate the required result when one of the components of the complex operand is infinite. (Explicit multiplication of the infinite component by the zero component obtained during promotion yields a NaN that propagates into the final result.) Analogous advice applies in the case of multiplication of a complex operand and a pure-imaginary operand, and in the case of division of a complex operand by a real or pure-imaginary operand. 56

Likewise, because the usual mathematical meaning of addition of a complex operand and a real operand is that the imaginary operand remains unchanged, an implementation should not perform this operation by first promoting the real operand to complex type and then performing a full complex addition. In implementations in which the Signed_Zeros attribute of the component type is True (and which therefore conform to IEC 559:1989 in regard to the handling of the sign of zero in predefined arithmetic operations), the latter technique will not generate the required result when the imaginary component of the complex operand is a negatively signed zero. (Explicit addition of the negative zero to the zero obtained during promotion yields a positive zero.) Analogous advice applies in the case of addition of a complex operand and a pure-imaginary operand, and in the case of subtraction of a complex operand and a real or pure-imaginary operand. 57

Implementations in which Real'Signed_Zeros is True should attempt to provide a rational treatment of the signs of zero results and result components. As one example, the result of the Argument function should have the sign of the imaginary component of the parameter X when the point represented by that 58

parameter lies on the positive real axis; as another, the sign of the imaginary component of the Compose_From_Polar function should be the same as (resp., the opposite of) that of the Argument parameter when that parameter has a value of zero and the Modulus parameter has a nonnegative (resp., negative) value.

G.1.2 Complex Elementary Functions

Static Semantics

- 1 The generic library package Numerics.Generic_Complex_Elementary_Functions has the following declaration:

```

2/2   with Ada.Numerics.Generic_Complex_Types;
generic
  with package Complex_Types is
    new Ada.Numerics.Generic_Complex_Types (<>);
  use Complex_Types;
package Ada.Numerics.Generic_Complex_Elementary_Functions is
  pragma Pure(Generic_Complex_Elementary_Functions);

3   function Sqrt (X : Complex)    return Complex;
  function Log  (X : Complex)    return Complex;
  function Exp  (X : Complex)    return Complex;
  function Exp  (X : Imaginary)  return Complex;
  function "***" (Left : Complex; Right : Complex)  return Complex;
  function "***" (Left : Complex; Right : Real'Base) return Complex;
  function "***" (Left : Real'Base; Right : Complex)  return Complex;

4   function Sin  (X : Complex)  return Complex;
  function Cos  (X : Complex)  return Complex;
  function Tan  (X : Complex)  return Complex;
  function Cot  (X : Complex)  return Complex;

5   function Arccsin (X : Complex) return Complex;
  function Arccos (X : Complex) return Complex;
  function Arctan (X : Complex) return Complex;
  function Arccot (X : Complex) return Complex;

6   function Sinh (X : Complex) return Complex;
  function Cosh (X : Complex) return Complex;
  function Tanh (X : Complex) return Complex;
  function Coth (X : Complex) return Complex;

7   function Arcsinh (X : Complex) return Complex;
  function Arccosh (X : Complex) return Complex;
  function Arctanh (X : Complex) return Complex;
  function Arccoth (X : Complex) return Complex;

8   end Ada.Numerics.Generic_Complex_Elementary_Functions;
```

- 9/1 The library package Numerics.Complex_Elementary_Functions is declared pure and defines the same subprograms as Numerics.Generic_Complex_Elementary_Functions, except that the predefined type Float is systematically substituted for Real'Base, and the Complex and Imaginary types exported by Numerics.Complex_Types are systematically substituted for Complex and Imaginary, throughout. Nongeneric equivalents of Numerics.Generic_Complex_Elementary_Functions corresponding to each of the other predefined floating point types are defined similarly, with the names Numerics.Short_Complex_Elementary_Functions, Numerics.Long_Complex_Elementary_Functions, etc.

- 10 The overloading of the Exp function for the pure-imaginary type is provided to give the user an alternate way to compose a complex value from a given modulus and argument. In addition to Compose_From_Polar(Rho, Theta) (see G.1.1), the programmer may write Rho * Exp(i * Theta).

- 11 The imaginary (resp., real) component of the parameter X of the forward hyperbolic (resp., trigonometric) functions and of the Exp function (and the parameter X, itself, in the case of the overloading of the Exp

function for the pure-imaginary type) represents an angle measured in radians, as does the imaginary (resp., real) component of the result of the Log and inverse hyperbolic (resp., trigonometric) functions.

The functions have their usual mathematical meanings. However, the arbitrariness inherent in the placement of branch cuts, across which some of the complex elementary functions exhibit discontinuities, is eliminated by the following conventions:

- The imaginary component of the result of the Sqrt and Log functions is discontinuous as the parameter X crosses the negative real axis. 13
- The result of the exponentiation operator when the left operand is of complex type is discontinuous as that operand crosses the negative real axis. 14
- The imaginary component of the result of the Arcsin, Arccos, and Arctanh functions is discontinuous as the parameter X crosses the real axis to the left of -1.0 or the right of 1.0 . 15/2
- The real component of the result of the Arctan and Arcsinh functions is discontinuous as the parameter X crosses the imaginary axis below $-i$ or above i . 16/2
- The real component of the result of the Arccot function is discontinuous as the parameter X crosses the imaginary axis below $-i$ or above i . 17/2
- The imaginary component of the Arccosh function is discontinuous as the parameter X crosses the real axis to the left of 1.0 . 18
- The imaginary component of the Arccoth function is discontinuous as the parameter X crosses the real axis between -1.0 and 1.0 . 19

The computed results of the mathematically multivalued functions are rendered single-valued by the following conventions, which are meant to imply that the principal branch is an analytic continuation of the corresponding real-valued function in Numerics.Generic_Elementary_Functions. (For Arctan and Arcot, the single-argument function in question is that obtained from the two-argument version by fixing the second argument to be its default value.) 20/2

- The real component of the result of the Sqrt and Arccosh functions is nonnegative. 21
- The same convention applies to the imaginary component of the result of the Log function as applies to the result of the natural-cycle version of the Argument function of Numerics.Generic_Complex_Types (see G.1.1). 22
- The range of the real (resp., imaginary) component of the result of the Arcsin and Arctan (resp., Arcsinh and Arctanh) functions is approximately $-\pi/2.0$ to $\pi/2.0$. 23
- The real (resp., imaginary) component of the result of the Arccos and Arcot (resp., Arccoth) functions ranges from 0.0 to approximately π . 24
- The range of the imaginary component of the result of the Arccosh function is approximately $-\pi$ to π . 25

In addition, the exponentiation operator inherits the single-valuedness of the Log function. 26

Dynamic Semantics

The exception Numerics.Argument_Error is raised by the exponentiation operator, signaling a parameter value outside the domain of the corresponding mathematical function, when the value of the left operand is zero and the real component of the exponent (or the exponent itself, when it is of real type) is zero. 27

The exception Constraint_Error is raised, signaling a pole of the mathematical function (analogous to dividing by zero), in the following cases, provided that Complex_Types.Real'Machine_Overflows is True: 28

- by the Log, Cot, and Coth functions, when the value of the parameter X is zero; 29

- 30 • by the exponentiation operator, when the value of the left operand is zero and the real
31 component of the exponent (or the exponent itself, when it is of real type) is negative;
32 • by the Arctan and Arccot functions, when the value of the parameter X is $\pm i$;
33 • by the Arctanh and Arccoth functions, when the value of the parameter X is ± 1.0 .
34 Constraint_Error can also be raised when a finite result overflows (see G.2.6); this may occur for
35 parameter values sufficiently *near* poles, and, in the case of some of the functions, for parameter values
36 having components of sufficiently large magnitude. When Complex_Types.Real'Machine_Overflows is
37 False, the result at poles is unspecified.

Implementation Requirements

- 34 In the implementation of Numerics.Generic_Complex_Elementary_Functions, the range of intermediate
35 values allowed during the calculation of a final result shall not be affected by any range constraint of the
36 subtype Complex_Types.Real.
37 In the following cases, evaluation of a complex elementary function shall yield the *prescribed result* (or a
38 result having the prescribed component), provided that the preceding rules do not call for an exception to
39 be raised:
40 • When the parameter X has the value zero, the Sqrt, Sin, Arcsin, Tan, Arctan, Sinh, Arcsinh,
41 Tanh, and Arctanh functions yield a result of zero; the Exp, Cos, and Cosh functions yield a
42 result of one; the Arccos and Arccot functions yield a real result; and the Arccoth function yields
43 an imaginary result.
44 • When the parameter X has the value one, the Sqrt function yields a result of one; the Log,
45 Arccos, and Arccosh functions yield a result of zero; and the Arcsin function yields a real result.
46 • When the parameter X has the value -1.0 , the Sqrt function yields the result
47 • i (resp., $-i$), when the sign of the imaginary component of X is positive (resp., negative), if
48 Complex_Types.Real'Signed_Zeros is True;
49 • i , if Complex_Types.Real'Signed_Zeros is False;
50 • When the parameter X has the value -1.0 , the Log function yields an imaginary result; and the
51 Arcsin and Arccos functions yield a real result.
52 • When the parameter X has the value $\pm i$, the Log function yields an imaginary result.
53 • Exponentiation by a zero exponent yields the value one. Exponentiation by a unit exponent
54 yields the value of the left operand (as a complex value). Exponentiation of the value one yields
55 the value one. Exponentiation of the value zero yields the value zero.
56 Other accuracy requirements for the complex elementary functions, which apply only in the strict mode,
57 are given in G.2.6.
58 The sign of a zero result or zero result component yielded by a complex elementary function is
59 implementation defined when Complex_Types.Real'Signed_Zeros is True.

Implementation Permissions

- 60 The nongeneric equivalent packages may, but need not, be actual instantiations of the generic package
61 with the appropriate predefined nongeneric equivalent of Numerics.Generic_Complex_Types; if they are,
62 then the latter shall have been obtained by actual instantiation of Numerics.Generic_Complex_Types.

The exponentiation operator may be implemented in terms of the Exp and Log functions. Because this implementation yields poor accuracy in some parts of the domain, no accuracy requirement is imposed on complex exponentiation.

The implementation of the Exp function of a complex parameter X is allowed to raise the exception Constraint_Error, signaling overflow, when the real component of X exceeds an unspecified threshold that is approximately log(Complex_Types.Real'Safe_Last). This permission recognizes the impracticality of avoiding overflow in the marginal case that the exponential of the real component of X exceeds the safe range of Complex_Types.Real but both components of the final result do not. Similarly, the Sin and Cos (resp., Sinh and Cosh) functions are allowed to raise the exception Constraint_Error, signaling overflow, when the absolute value of the imaginary (resp., real) component of the parameter X exceeds an unspecified threshold that is approximately log(Complex_Types.Real'Safe_Last) + log(2.0). This permission recognizes the impracticality of avoiding overflow in the marginal case that the hyperbolic sine or cosine of the imaginary (resp., real) component of X exceeds the safe range of Complex_Types.Real but both components of the final result do not.

Implementation Advice

Implementations in which Complex_Types.Real'Signed_Zeros is True should attempt to provide a rational treatment of the signs of zero results and result components. For example, many of the complex elementary functions have components that are odd functions of one of the parameter components; in these cases, the result component should have the sign of the parameter component at the origin. Other complex elementary functions have zero components whose sign is opposite that of a parameter component at the origin, or is always positive or always negative.

G.1.3 Complex Input-Output

The generic package Text_IO.Complex_IO defines procedures for the formatted input and output of complex values. The generic actual parameter in an instantiation of Text_IO.Complex_IO is an instance of Numerics.Generic_Complex_Types for some floating point subtype. Exceptional conditions are reported by raising the appropriate exception defined in Text_IO.

Static Semantics

The generic library package Text_IO.Complex_IO has the following declaration:

```

with Ada.Numerics.Generic_Complex_Types;
generic
  with package Complex_Types is
    new Ada.Numerics.Generic_Complex_Types (<>);
package Ada.Text_IO.Complex_IO is
  use Complex_Types;
  Default_Fore : Field := 2;
  Default_Aft  : Field := Real'Digits - 1;
  Default_Exp   : Field := 3;
  procedure Get (File   : in File_Type;
                 Item   : out Complex;
                 Width  : in Field := 0);
  procedure Get (Item   : out Complex;
                 Width  : in Field := 0);

```

```

7      procedure Put (File : in File_Type;
                      Item : in Complex;
                      Fore : in Field := Default_Fore;
                      Aft : in Field := Default_Aft;
                      Exp : in Field := Default_Exp);
8      procedure Put (Item : in Complex;
                      Fore : in Field := Default_Fore;
                      Aft : in Field := Default_Aft;
                      Exp : in Field := Default_Exp);
9      procedure Get (From : in String;
                      Item : out Complex;
                      Last : out Positive);
10     procedure Put (To : out String;
                      Item : in Complex;
                      Aft : in Field := Default_Aft;
                      Exp : in Field := Default_Exp);

11    end Ada.Text_IO.Complex_IO;

```

9.1/2 The library package Complex_Text_IO defines the same subprograms as Text_IO.Complex_IO, except that the predefined type Float is systematically substituted for Real, and the type Numerics.Complex_Types.Complex is systematically substituted for Complex throughout. Non-generic equivalents of Text_IO.Complex_IO corresponding to each of the other predefined floating point types are defined similarly, with the names Short_Complex_Text_IO, Long_Complex_Text_IO, etc.

10 The semantics of the Get and Put procedures are as follows:

```

11  procedure Get (File : in File_Type;
                  Item : out Complex;
                  Width : in Field := 0);
12  procedure Get (Item : out Complex;
                  Width : in Field := 0);

```

12/1 The input sequence is a pair of optionally signed real literals representing the real and imaginary components of a complex value. These components have the format defined for the corresponding Get procedure of an instance of Text_IO.Float_IO (see A.10.9) for the base subtype of Complex_Types.Real. The pair of components may be separated by a comma or surrounded by a pair of parentheses or both. Blanks are freely allowed before each of the components and before the parentheses and comma, if either is used. If the value of the parameter Width is zero, then

- 13 • line and page terminators are also allowed in these places;
- 14 • the components shall be separated by at least one blank or line terminator if the comma is omitted; and
- 15 • reading stops when the right parenthesis has been read, if the input sequence includes a left parenthesis, or when the imaginary component has been read, otherwise.

15.1 If a nonzero value of Width is supplied, then

- 16 • the components shall be separated by at least one blank if the comma is omitted; and
- 17 • exactly Width characters are read, or the characters (possibly none) up to a line terminator, whichever comes first (blanks are included in the count).

18 Returns, in the parameter Item, the value of type Complex that corresponds to the input sequence.

19 The exception Text_IO.Data_Error is raised if the input sequence does not have the required syntax or if the components of the complex value obtained are not of the base subtype of Complex_Types.Real.

```
procedure Put (File : in File_Type;
               Item : in Complex;
               Fore : in Field := Default_Fore;
               Aft : in Field := Default_Aft;
               Exp : in Field := Default_Exp);
procedure Put (Item : in Complex;
               Fore : in Field := Default_Fore;
               Aft : in Field := Default_Aft;
               Exp : in Field := Default_Exp);
```

Outputs the value of the parameter Item as a pair of decimal literals representing the real and imaginary components of the complex value, using the syntax of an aggregate. More specifically,

- outputs a left parenthesis;
- outputs the value of the real component of the parameter Item with the format defined by the corresponding Put procedure of an instance of Text_IO.Float_IO for the base subtype of Complex_Types.Real, using the given values of Fore, Aft, and Exp;
- outputs a comma;
- outputs the value of the imaginary component of the parameter Item with the format defined by the corresponding Put procedure of an instance of Text_IO.Float_IO for the base subtype of Complex_Types.Real, using the given values of Fore, Aft, and Exp;
- outputs a right parenthesis.

```
procedure Get (From : in String;
               Item : out Complex;
               Last : out Positive);
```

Reads a complex value from the beginning of the given string, following the same rule as the Get procedure that reads a complex value from a file, but treating the end of the string as a file terminator. Returns, in the parameter Item, the value of type Complex that corresponds to the input sequence. Returns in Last the index value such that From(Last) is the last character read.

The exception Text_IO.Data_Error is raised if the input sequence does not have the required syntax or if the components of the complex value obtained are not of the base subtype of Complex_Types.Real.

```
procedure Put (To : out String;
               Item : in Complex;
               Aft : in Field := Default_Aft;
               Exp : in Field := Default_Exp);
```

Outputs the value of the parameter Item to the given string as a pair of decimal literals representing the real and imaginary components of the complex value, using the syntax of an aggregate. More specifically,

- a left parenthesis, the real component, and a comma are left justified in the given string, with the real component having the format defined by the Put procedure (for output to a file) of an instance of Text_IO.Float_IO for the base subtype of Complex_Types.Real, using a value of zero for Fore and the given values of Aft and Exp;
- the imaginary component and a right parenthesis are right justified in the given string, with the imaginary component having the format defined by the Put procedure (for output to a file) of an instance of Text_IO.Float_IO for the base subtype of Complex_Types.Real, using a value for Fore that completely fills the remainder of the string, together with the given values of Aft and Exp.

- 34 The exception `Text_IO.Layout_Error` is raised if the given string is too short to hold the formatted output.

Implementation Permissions

- 35 Other exceptions declared (by renaming) in `Text_IO` may be raised by the preceding procedures in the appropriate circumstances, as for the corresponding procedures of `Text_IO.Float_IO`.

G.1.4 The Package `Wide_Text_IO.Complex_IO`

Static Semantics

- 1 Implementations shall also provide the generic library package `Wide_Text_IO.Complex_IO`. Its declaration is obtained from that of `Text_IO.Complex_IO` by systematically replacing `Text_IO` by `Wide_Text_IO` and `String` by `Wide_String`; the description of its behavior is obtained by additionally replacing references to particular characters (commas, parentheses, etc.) by those for the corresponding wide characters.

G.1.5 The Package `Wide_Wide_Text_IO.Complex_IO`

Static Semantics

- 1/2 Implementations shall also provide the generic library package `Wide_Wide_Text_IO.Complex_IO`. Its declaration is obtained from that of `Text_IO.Complex_IO` by systematically replacing `Text_IO` by `Wide_Wide_Text_IO` and `String` by `Wide_Wide_String`; the description of its behavior is obtained by additionally replacing references to particular characters (commas, parentheses, etc.) by those for the corresponding wide wide characters.

G.2 Numeric Performance Requirements

Implementation Requirements

- 1 Implementations shall provide a user-selectable mode in which the accuracy and other numeric performance requirements detailed in the following subclauses are observed. This mode, referred to as the *strict mode*, may or may not be the default mode; it directly affects the results of the predefined arithmetic operations of real types and the results of the subprograms in children of the `Numerics` package, and indirectly affects the operations in other language defined packages. Implementations shall also provide the opposing mode, which is known as the *relaxed mode*.

Implementation Permissions

- 2 Either mode may be the default mode.
- 3 The two modes need not actually be different.

G.2.1 Model of Floating Point Arithmetic

In the strict mode, the predefined operations of a floating point type shall satisfy the accuracy requirements specified here and shall avoid or signal overflow in the situations described. This behavior is presented in terms of a model of floating point arithmetic that builds on the concept of the canonical form (see A.5.3).

Static Semantics

Associated with each floating point type is an infinite set of model numbers. The model numbers of a type are used to define the accuracy requirements that have to be satisfied by certain predefined operations of the type; through certain attributes of the model numbers, they are also used to explain the meaning of a user-declared floating point type declaration. The model numbers of a derived type are those of the parent type; the model numbers of a subtype are those of its type.

The *model numbers* of a floating point type T are zero and all the values expressible in the canonical form (for the type T), in which *mantissa* has T'Model_Mantissa digits and *exponent* has a value greater than or equal to T'Model_Emin. (These attributes are defined in G.2.2.)

A *model interval* of a floating point type is any interval whose bounds are model numbers of the type. The *model interval* of a type T associated with a value v is the smallest model interval of T that includes v. (The model interval associated with a model number of a type consists of that number only.)

Implementation Requirements

The accuracy requirements for the evaluation of certain predefined operations of floating point types are as follows.

An *operand interval* is the model interval, of the type specified for the operand of an operation, associated with the value of the operand.

For any predefined arithmetic operation that yields a result of a floating point type T, the required bounds on the result are given by a model interval of T (called the *result interval*) defined in terms of the operand values as follows:

- The result interval is the smallest model interval of T that includes the minimum and the maximum of all the values obtained by applying the (exact) mathematical operation to values arbitrarily selected from the respective operand intervals.

The result interval of an exponentiation is obtained by applying the above rule to the sequence of multiplications defined by the exponent, assuming arbitrary association of the factors, and to the final division in the case of a negative exponent.

The result interval of a conversion of a numeric value to a floating point type T is the model interval of T associated with the operand value, except when the source expression is of a fixed point type with a *small* that is not a power of T'Machine_Radix or is a fixed point multiplication or division either of whose operands has a *small* that is not a power of T'Machine_Radix; in these cases, the result interval is implementation defined.

For any of the foregoing operations, the implementation shall deliver a value that belongs to the result interval when both bounds of the result interval are in the safe range of the result type T, as determined by the values of T'Safe_First and T'Safe_Last; otherwise,

- if T'Machine_Overflows is True, the implementation shall either deliver a value that belongs to the result interval or raise Constraint_Error;

- 13 • if $T\text{Machine_Overflows}$ is False, the result is implementation defined.
- 14 For any predefined relation on operands of a floating point type T , the implementation may deliver any value (i.e., either True or False) obtained by applying the (exact) mathematical comparison to values arbitrarily chosen from the respective operand intervals.
- 15 The result of a membership test is defined in terms of comparisons of the operand value with the lower and upper bounds of the given range or type mark (the usual rules apply to these comparisons).

Implementation Permissions

- 16 If the underlying floating point hardware implements division as multiplication by a reciprocal, the result interval for division (and exponentiation by a negative exponent) is implementation defined.

G.2.2 Model-Oriented Attributes of Floating Point Types

- 1 In implementations that support the Numerics Annex, the model-oriented attributes of floating point types shall yield the values defined here, in both the strict and the relaxed modes. These definitions add conditions to those in A.5.3.

Static Semantics

- 2 For every subtype S of a floating point type T :

3/2 $S\text{Model_Mantissa}$

Yields the number of digits in the mantissa of the canonical form of the model numbers of T (see A.5.3). The value of this attribute shall be greater than or equal to

$$\lceil d \cdot \log(10) / \log(T\text{Machine_Radix}) \rceil + g$$

3.2/2 where d is the requested decimal precision of T , and g is 0 if $T\text{Machine_Radix}$ is a positive power of 10 and 1 otherwise. In addition, $T\text{Model_Mantissa}$ shall be less than or equal to the value of $T\text{Machine_Mantissa}$. This attribute yields a value of the type *universal_integer*.

4 $S\text{Model_Emin}$

Yields the minimum exponent of the canonical form of the model numbers of T (see A.5.3). The value of this attribute shall be greater than or equal to the value of $T\text{Machine_Emin}$. This attribute yields a value of the type *universal_integer*.

5 $S\text{Safe_First}$

Yields the lower bound of the safe range of T . The value of this attribute shall be a model number of T and greater than or equal to the lower bound of the base range of T . In addition, if T is declared by a *floating_point_definition* or is derived from such a type, and the *floating_point_definition* includes a *real_range_specification* specifying a lower bound of lb , then the value of this attribute shall be less than or equal to lb ; otherwise, it shall be less than or equal to -10.0^{4-d} , where d is the requested decimal precision of T . This attribute yields a value of the type *universal_real*.

6 $S\text{Safe_Last}$

Yields the upper bound of the safe range of T . The value of this attribute shall be a model number of T and less than or equal to the upper bound of the base range of T . In addition, if T is declared by a *floating_point_definition* or is derived from such a type, and the *floating_point_definition* includes a *real_range_specification* specifying an upper bound of ub , then the value of this attribute shall be greater than or equal to ub ; otherwise, it shall be greater than or equal to 10.0^{4-d} , where d is the requested decimal precision of T . This attribute yields a value of the type *universal_real*.

S'Model Denotes a function (of a parameter X) whose specification is given in A.5.3. If X is a model number of T , the function yields X ; otherwise, it yields the value obtained by rounding or truncating X to either one of the adjacent model numbers of T . Constraint_Error is raised if the resulting model number is outside the safe range of S. A zero result has the sign of X when S'Signed_Zeros is True.

Subject to the constraints given above, the values of S'Model_Mantissa and S'Safe_Last are to be maximized, and the values of S'Model_Emin and S'Safe_First minimized, by the implementation as follows:

- First, S'Model_Mantissa is set to the largest value for which values of S'Model_Emin, S'Safe_First, and S'Safe_Last can be chosen so that the implementation satisfies the strict-mode requirements of G.2.1 in terms of the model numbers and safe range induced by these attributes.
- Next, S'Model_Emin is set to the smallest value for which values of S'Safe_First and S'Safe_Last can be chosen so that the implementation satisfies the strict-mode requirements of G.2.1 in terms of the model numbers and safe range induced by these attributes and the previously determined value of S'Model_Mantissa.
- Finally, S'Safe_First and S'Safe_Last are set (in either order) to the smallest and largest values, respectively, for which the implementation satisfies the strict-mode requirements of G.2.1 in terms of the model numbers and safe range induced by these attributes and the previously determined values of S'Model_Mantissa and S'Model_Emin.

G.2.3 Model of Fixed Point Arithmetic

In the strict mode, the predefined arithmetic operations of a fixed point type shall satisfy the accuracy requirements specified here and shall avoid or signal overflow in the situations described.

Implementation Requirements

The accuracy requirements for the predefined fixed point arithmetic operations and conversions, and the results of relations on fixed point operands, are given below.

The operands of the fixed point adding operators, absolute value, and comparisons have the same type. These operations are required to yield exact results, unless they overflow.

Multiplications and divisions are allowed between operands of any two fixed point types; the result has to be (implicitly or explicitly) converted to some other numeric type. For purposes of defining the accuracy rules, the multiplication or division and the conversion are treated as a single operation whose accuracy depends on three types (those of the operands and the result). For decimal fixed point types, the attribute T'Round may be used to imply explicit conversion with rounding (see 3.5.10).

When the result type is a floating point type, the accuracy is as given in G.2.1. For some combinations of the operand and result types in the remaining cases, the result is required to belong to a small set of values called the *perfect result set*; for other combinations, it is required merely to belong to a generally larger and implementation-defined set of values called the *close result set*. When the result type is a decimal fixed point type, the perfect result set contains a single value; thus, operations on decimal types are always fully specified.

When one operand of a fixed-fixed multiplication or division is of type *universal_real*, that operand is not implicitly converted in the usual sense, since the context does not determine a unique target type, but the accuracy of the result of the multiplication or division (i.e., whether the result has to belong to the perfect result set or merely the close result set) depends on the value of the operand of type *universal_real* and on the types of the other operand and of the result.

7 For a fixed point multiplication or division whose (exact) mathematical result is v , and for the conversion
of a value v to a fixed point type, the perfect result set and close result set are defined as follows:

- 8 • If the result type is an ordinary fixed point type with a *small* of s ,
 - 9 • if v is an integer multiple of s , then the perfect result set contains only the value v ;
 - 10 • otherwise, it contains the integer multiple of s just below v and the integer multiple of s just above v .

11 The close result set is an implementation-defined set of consecutive integer multiples of s containing the perfect result set as a subset.

- 12 • If the result type is a decimal type with a *small* of s ,
 - 13 • if v is an integer multiple of s , then the perfect result set contains only the value v ;
 - 14 • otherwise, if truncation applies then it contains only the integer multiple of s in the direction toward zero, whereas if rounding applies then it contains only the nearest integer multiple of s (with ties broken by rounding away from zero).

15 The close result set is an implementation-defined set of consecutive integer multiples of s containing the perfect result set as a subset.

- 16 • If the result type is an integer type,
 - 17 • if v is an integer, then the perfect result set contains only the value v ;
 - 18 • otherwise, it contains the integer nearest to the value v (if v lies equally distant from two consecutive integers, the perfect result set contains the one that is further from zero).

19 The close result set is an implementation-defined set of consecutive integers containing the perfect result set as a subset.

20 The result of a fixed point multiplication or division shall belong either to the perfect result set or to the close result set, as described below, if overflow does not occur. In the following cases, if the result type is a fixed point type, let s be its *small*; otherwise, i.e. when the result type is an integer type, let s be 1.0.

- 21 • For a multiplication or division neither of whose operands is of type *universal_real*, let l and r be the *smalls* of the left and right operands. For a multiplication, if $(l \cdot r) / s$ is an integer or the reciprocal of an integer (the *smalls* are said to be “compatible” in this case), the result shall belong to the perfect result set; otherwise, it belongs to the close result set. For a division, if $l / (r \cdot s)$ is an integer or the reciprocal of an integer (i.e., the *smalls* are compatible), the result shall belong to the perfect result set; otherwise, it belongs to the close result set.
- 22 • For a multiplication or division having one *universal_real* operand with a value of v , note that it is always possible to factor v as an integer multiple of a “compatible” *small*, but the integer multiple may be “too big.” If there exists a factorization in which that multiple is less than some implementation-defined limit, the result shall belong to the perfect result set; otherwise, it belongs to the close result set.

23 A multiplication $P * Q$ of an operand of a fixed point type F by an operand of an integer type I , or vice-versa, and a division P / Q of an operand of a fixed point type F by an operand of an integer type I , are also allowed. In these cases, the result has a type of F ; explicit conversion of the result is never required. The accuracy required in these cases is the same as that required for a multiplication $F(P * Q)$ or a division $F(P / Q)$ obtained by interpreting the operand of the integer type to have a fixed point type with a *small* of 1.0.

24 The accuracy of the result of a conversion from an integer or fixed point type to a fixed point type, or from a fixed point type to an integer type, is the same as that of a fixed point multiplication of the source value by a fixed point operand having a *small* of 1.0 and a value of 1.0, as given by the foregoing rules. The result of a conversion from a floating point type to a fixed point type shall belong to the close result set.

The result of a conversion of a *universal_real* operand to a fixed point type shall belong to the perfect result set.

The possibility of overflow in the result of a predefined arithmetic operation or conversion yielding a result of a fixed point type T is analogous to that for floating point types, except for being related to the base range instead of the safe range. If all of the permitted results belong to the base range of T, then the implementation shall deliver one of the permitted results; otherwise,

- if T'Machine_Overflows is True, the implementation shall either deliver one of the permitted results or raise Constraint_Error;
- if T'Machine_Overflows is False, the result is implementation defined.

G.2.4 Accuracy Requirements for the Elementary Functions

In the strict mode, the performance of Numerics.Generic_Elementary_Functions shall be as specified here.

Implementation Requirements

When an exception is not raised, the result of evaluating a function in an instance EF of Numerics.Generic_Elementary_Functions belongs to a *result interval*, defined as the smallest model interval of EF.Float_Type that contains all the values of the form $f \cdot (1.0 + d)$, where f is the exact value of the corresponding mathematical function at the given parameter values, d is a real number, and |d| is less than or equal to the function's *maximum relative error*. The function delivers a value that belongs to the result interval when both of its bounds belong to the safe range of EF.Float_Type; otherwise,

- if EF.Float_Type'Machine_Overflows is True, the function either delivers a value that belongs to the result interval or raises Constraint_Error, signaling overflow;
- if EF.Float_Type'Machine_Overflows is False, the result is implementation defined.

The maximum relative error exhibited by each function is as follows:

- $2.0 \cdot EF.Float_Type'Model_Epsilon$, in the case of the Sqrt, Sin, and Cos functions;
- $4.0 \cdot EF.Float_Type'Model_Epsilon$, in the case of the Log, Exp, Tan, Cot, and inverse trigonometric functions; and
- $8.0 \cdot EF.Float_Type'Model_Epsilon$, in the case of the forward and inverse hyperbolic functions.

The maximum relative error exhibited by the exponentiation operator, which depends on the values of the operands, is $(4.0 + |\text{Right} \cdot \log(\text{Left})| / 32.0) \cdot EF.Float_Type'Model_Epsilon$.

The maximum relative error given above applies throughout the domain of the forward trigonometric functions when the Cycle parameter is specified. When the Cycle parameter is omitted, the maximum relative error given above applies only when the absolute value of the angle parameter X is less than or equal to some implementation-defined *angle threshold*, which shall be at least EF.Float_Type'Machine_Radix^{└EF.Float_Type'Machine_Mantissa/2┘}. Beyond the angle threshold, the accuracy of the forward trigonometric functions is implementation defined.

The prescribed results specified in A.5.1 for certain functions at particular parameter values take precedence over the maximum relative error bounds; effectively, they narrow to a single value the result interval allowed by the maximum relative error bounds. Additional rules with a similar effect are given by table G-1 for the inverse trigonometric functions, at particular parameter values for which the mathematical result is possibly not a model number of EF.Float_Type (or is, indeed, even transcendental). In each table entry, the values of the parameters are such that the result lies on the axis between two

quadrants; the corresponding accuracy rule, which takes precedence over the maximum relative error bounds, is that the result interval is the model interval of *EF*.*Float_Type* associated with the exact mathematical result given in the table.

12/1 This paragraph was deleted.

13 The last line of the table is meant to apply when *EF*.*Float_Type*'Signed_Zeros is False; the two lines just above it, when *EF*.*Float_Type*'Signed_Zeros is True and the parameter Y has a zero value with the indicated sign.

Table G-1: Tightly Approximated Elementary Function Results

Function	Value of X	Value of Y	Exact Result when Cycle Specified	Exact Result when Cycle Omitted
Arcsin	1.0	n.a.	Cycle/4.0	$\pi/2.0$
Arcsin	-1.0	n.a.	-Cycle/4.0	$-\pi/2.0$
Arccos	0.0	n.a.	Cycle/4.0	$\pi/2.0$
Arccos	-1.0	n.a.	Cycle/2.0	π
Arctan and Arccot	0.0	positive	Cycle/4.0	$\pi/2.0$
Arctan and Arccot	0.0	negative	-Cycle/4.0	$-\pi/2.0$
Arctan and Arccot	negative	+0.0	Cycle/2.0	π
Arctan and Arccot	negative	-0.0	-Cycle/2.0	$-\pi$
Arctan and Arccot	negative	0.0	Cycle/2.0	π

14 The amount by which the result of an inverse trigonometric function is allowed to spill over into a quadrant adjacent to the one corresponding to the principal branch, as given in A.5.1, is limited. The rule is that the result belongs to the smallest model interval of *EF*.*Float_Type* that contains both boundaries of the quadrant corresponding to the principal branch. This rule also takes precedence over the maximum relative error bounds, effectively narrowing the result interval allowed by them.

15 Finally, the following specifications also take precedence over the maximum relative error bounds:

- 16 • The absolute value of the result of the Sin, Cos, and Tanh functions never exceeds one.
- 17 • The absolute value of the result of the Coth function is never less than one.
- 18 • The result of the Cosh function is never less than one.

Implementation Advice

19 The versions of the forward trigonometric functions without a Cycle parameter should not be implemented by calling the corresponding version with a Cycle parameter of 2.0*Numerics.Pi, since this will not provide the required accuracy in some portions of the domain. For the same reason, the version of Log without a Base parameter should not be implemented by calling the corresponding version with a Base parameter of Numerics.e.

G.2.5 Performance Requirements for Random Number Generation

In the strict mode, the performance of Numerics.Float_Random and Numerics.Discrete_Random shall be as specified here.

Implementation Requirements

Two different calls to the time-dependent Reset procedure shall reset the generator to different states, provided that the calls are separated in time by at least one second and not more than fifty years.

The implementation's representations of generator states and its algorithms for generating random numbers shall yield a period of at least $2^{31}-2$; much longer periods are desirable but not required.

The implementations of Numerics.Float_Random.Random and Numerics.Discrete_Random.Random shall pass at least 85% of the individual trials in a suite of statistical tests. For Numerics.Float_Random, the tests are applied directly to the floating point values generated (i.e., they are not converted to integers first), while for Numerics.Discrete_Random they are applied to the generated values of various discrete types. Each test suite performs 6 different tests, with each test repeated 10 times, yielding a total of 60 individual trials. An individual trial is deemed to pass if the chi-square value (or other statistic) calculated for the observed counts or distribution falls within the range of values corresponding to the 2.5 and 97.5 percentage points for the relevant degrees of freedom (i.e., it shall be neither too high nor too low). For the purpose of determining the degrees of freedom, measurement categories are combined whenever the expected counts are fewer than 5.

G.2.6 Accuracy Requirements for Complex Arithmetic

In the strict mode, the performance of Numerics.Generic_Complex_Types and Numerics.Generic_Complex_Elementary_Functions shall be as specified here.

Implementation Requirements

When an exception is not raised, the result of evaluating a real function of an instance CT of Numerics.Generic_Complex_Types (i.e., a function that yields a value of subtype $CT.\text{Real}'\text{Base}$ or $CT.\text{Imaginary}$) belongs to a result interval defined as for a real elementary function (see G.2.4).

When an exception is not raised, each component of the result of evaluating a complex function of such an instance, or of an instance of Numerics.Generic_Complex_Elementary_Functions obtained by instantiating the latter with CT (i.e., a function that yields a value of subtype $CT.\text{Complex}$), also belongs to a *result interval*. The result intervals for the components of the result are either defined by a *maximum relative error* bound or by a *maximum box error* bound. When the result interval for the real (resp., imaginary) component is defined by maximum relative error, it is defined as for that of a real function, relative to the exact value of the real (resp., imaginary) part of the result of the corresponding mathematical function. When defined by maximum box error, the result interval for a component of the result is the smallest model interval of $CT.\text{Real}$ that contains all the values of the corresponding part of $f \cdot (1.0 + d)$, where f is the exact complex value of the corresponding mathematical function at the given parameter values, d is complex, and $|d|$ is less than or equal to the given maximum box error. The function delivers a value that belongs to the result interval (or a value both of whose components belong to their respective result intervals) when both bounds of the result interval(s) belong to the safe range of $CT.\text{Real}$; otherwise,

- 4 • if *CT.Real'Machine_Overflows* is True, the function either delivers a value that belongs to the result interval (or a value both of whose components belong to their respective result intervals) or raises *Constraint_Error*, signaling overflow;
- 5 • if *CT.Real'Machine_Overflows* is False, the result is implementation defined.
- 6/2 The error bounds for particular complex functions are tabulated in table G-2. In the table, the error bound is given as the coefficient of *CT.Real'Model_Epsilon*.
- 7/1 *This paragraph was deleted.*

Table G-2: Error Bounds for Particular Complex Functions			
Function or Operator	Nature of Result	Nature of Bound	Error Bound
Modulus	real	max. rel. error	3.0
Argument	real	max. rel. error	4.0
Compose_From_Polar	complex	max. rel. error	3.0
"*" (both operands complex)	complex	max. box error	5.0
"/" (right operand complex)	complex	max. box error	13.0
Sqrt	complex	max. rel. error	6.0
Log	complex	max. box error	13.0
Exp (complex parameter)	complex	max. rel. error	7.0
Exp (imaginary parameter)	complex	max. rel. error	2.0
Sin, Cos, Sinh, and Cosh	complex	max. rel. error	11.0
Tan, Cot, Tanh, and Coth	complex	max. rel. error	35.0
inverse trigonometric	complex	max. rel. error	14.0
inverse hyperbolic	complex	max. rel. error	14.0

- 8 The maximum relative error given above applies throughout the domain of the *Compose_From_Polar* function when the *Cycle* parameter is specified. When the *Cycle* parameter is omitted, the maximum relative error applies only when the absolute value of the parameter *Argument* is less than or equal to the angle threshold (see G.2.4). For the *Exp* function, and for the forward hyperbolic (resp., trigonometric) functions, the maximum relative error given above likewise applies only when the absolute value of the imaginary (resp., real) component of the parameter *X* (or the absolute value of the parameter itself, in the case of the *Exp* function with a parameter of pure-imaginary type) is less than or equal to the angle threshold. For larger angles, the accuracy is implementation defined.
- 9 The prescribed results specified in G.1.2 for certain functions at particular parameter values take precedence over the error bounds; effectively, they narrow to a single value the result interval allowed by the error bounds for a component of the result. Additional rules with a similar effect are given below for certain inverse trigonometric and inverse hyperbolic functions, at particular parameter values for which a component of the mathematical result is transcendental. In each case, the accuracy rule, which takes precedence over the error bounds, is that the result interval for the stated result component is the model interval of *CT.Real* associated with the component's exact mathematical value. The cases in question are as follows:

- When the parameter X has the value zero, the real (resp., imaginary) component of the result of the Arccot (resp., Arccoth) function is in the model interval of *CT.Real* associated with the value $\pi/2.0$. 10
- When the parameter X has the value one, the real component of the result of the Arcsin function is in the model interval of *CT.Real* associated with the value $\pi/2.0$. 11
- When the parameter X has the value -1.0 , the real component of the result of the Arcsin (resp., Arccos) function is in the model interval of *CT.Real* associated with the value $-\pi/2.0$ (resp., π). 12

The amount by which a component of the result of an inverse trigonometric or inverse hyperbolic function is allowed to spill over into a quadrant adjacent to the one corresponding to the principal branch, as given in G.1.2, is limited. The rule is that the result belongs to the smallest model interval of *CT.Real* that contains both boundaries of the quadrant corresponding to the principal branch. This rule also takes precedence over the maximum error bounds, effectively narrowing the result interval allowed by them. 13/2

Finally, the results allowed by the error bounds are narrowed by one further rule: The absolute value of each component of the result of the Exp function, for a pure-imaginary parameter, never exceeds one. 14

Implementation Advice

The version of the Compose_From_Polar function without a Cycle parameter should not be implemented by calling the corresponding version with a Cycle parameter of $2.0 * \text{Numerics.Pi}$, since this will not provide the required accuracy in some portions of the domain. 15

G.3 Vector and Matrix Manipulation

Types and operations for the manipulation of real vectors and matrices are provided in Generic_Real_Arrays, which is defined in G.3.1. Types and operations for the manipulation of complex vectors and matrices are provided in Generic_Complex_Arrays, which is defined in G.3.2. Both of these library units are generic children of the predefined package Numerics (see A.5). Nongeneric equivalents of these packages for each of the predefined floating point types are also provided as children of Numerics. 1/2

G.3.1 Real Vectors and Matrices

Static Semantics

The generic library package Numerics.Generic_Real_Arrays has the following declaration: 1/2

```
generic
  type Real is digits <>;
package Ada.Numerics.Generic_Real_Arrays is
  pragma Pure(Generic_Real_Arrays);
  -- Types
  type Real_Vector is array (Integer range <>) of Real'Base;
  type Real_Matrix is array (Integer range <>, Integer range <>)
    of Real'Base;
  -- Subprograms for Real_Vector types
  -- Real_Vector arithmetic operations
  function "+" (Right : Real_Vector)      return Real_Vector;
  function "-" (Right : Real_Vector)      return Real_Vector;
  function "abs" (Right : Real_Vector)     return Real_Vector;
  function "+" (Left, Right : Real_Vector) return Real_Vector;
  function "-" (Left, Right : Real_Vector) return Real_Vector;
  function "*" (Left, Right : Real_Vector) return Real'Base;
```

```

10/2      function "abs" (Right : Real_Vector)           return Real'Base;
11/2      -- Real_Vector scaling operations
12/2      function "*" (Left : Real'Base;    Right : Real_Vector)
13/2          return Real_Vector;
14/2      function "*" (Left : Real_Vector; Right : Real'Base)
15/2          return Real_Vector;
16/2      function "/" (Left : Real_Vector; Right : Real'Base)
17/2          return Real_Vector;
18/2      -- Other Real_Vector operations
19/2      function Unit_Vector (Index : Integer;
20/2                      Order : Positive;
21/2                      First : Integer := 1) return Real_Vector;
22/2      -- Subprograms for Real_Matrix types
23/2      -- Real_Matrix arithmetic operations
24/2      function "+"      (Right : Real_Matrix) return Real_Matrix;
25/2      function "-"      (Right : Real_Matrix) return Real_Matrix;
26/2      function "abs"    (Right : Real_Matrix) return Real_Matrix;
27/2      function Transpose (X      : Real_Matrix) return Real_Matrix;
28/2      function "+"      (Left, Right : Real_Matrix) return Real_Matrix;
29/2      function "-"      (Left, Right : Real_Matrix) return Real_Matrix;
30/2      function "*"      (Left, Right : Real_Matrix) return Real_Matrix;
31/2      function "*"      (Left : Real_Vector; Right : Real_Matrix)
32/2          return Real_Vector;
33/2      function "*"      (Left : Real_Matrix; Right : Real_Vector)
34/2          return Real_Vector;
35/2      -- Real_Matrix scaling operations
36/2      function "*"      (Left : Real'Base;    Right : Real_Matrix)
37/2          return Real_Matrix;
38/2      function "*"      (Left : Real_Matrix; Right : Real'Base)
39/2          return Real_Matrix;
40/2      function "/"      (Left : Real_Matrix; Right : Real'Base)
41/2          return Real_Matrix;
42/2      -- Real_Matrix inversion and related operations
43/2      function Solve (A : Real_Matrix; X : Real_Vector) return Real_Vector;
44/2      function Solve (A, X : Real_Matrix) return Real_Matrix;
45/2      function Inverse (A : Real_Matrix) return Real_Matrix;
46/2      function Determinant (A : Real_Matrix) return Real'Base;
47/2      -- Eigenvalues and vectors of a real symmetric matrix
48/2      function Eigenvalues (A : Real_Matrix) return Real_Vector;
49/2      procedure Eigensystem (A      : in  Real_Matrix;
50/2                           Values : out Real_Vector;
51/2                           Vectors : out Real_Matrix);
52/2      -- Other Real_Matrix operations
53/2      function Unit_Matrix (Order      : Positive;
54/2                           First_1, First_2 : Integer := 1)
55/2          return Real_Matrix;
56/2  end Ada.Numerics.Generic_Real_Arrays;

```

31/2 The library package Numerics.Real_Arrays is declared pure and defines the same types and subprograms as Numerics.Generic_Real_Arrays, except that the predefined type Float is systematically substituted for Real'Base throughout. Nongeneric equivalents for each of the other predefined floating point types are defined similarly, with the names Numerics.Short_Real_Arrays, Numerics.Long_Real_Arrays, etc.

32/2 Two types are defined and exported by Numerics.Generic_Real_Arrays. The composite type Real_Vector is provided to represent a vector with components of type Real; it is defined as an unconstrained, one-

dimensional array with an index of type Integer. The composite type Real_Matrix is provided to represent a matrix with components of type Real; it is defined as an unconstrained, two-dimensional array with indices of type Integer.

The effect of the various subprograms is as described below. In most cases the subprograms are described in terms of corresponding scalar operations of the type Real; any exception raised by those operations is propagated by the array operation. Moreover, the accuracy of the result for each individual component is as defined for the scalar operation unless stated otherwise. 33/2

In the case of those operations which are defined to *involve an inner product*, Constraint_Error may be raised if an intermediate result is outside the range of Real'Base even though the mathematical final result would not be. 34/2

```
function "+" (Right : Real_Vector) return Real_Vector;
function "-" (Right : Real_Vector) return Real_Vector;
function "abs" (Right : Real_Vector) return Real_Vector;
```

Each operation returns the result of applying the corresponding operation of the type Real to each component of Right. The index range of the result is Right'Range. 36/2

```
function "+" (Left, Right : Real_Vector) return Real_Vector;
function "-" (Left, Right : Real_Vector) return Real_Vector;
```

Each operation returns the result of applying the corresponding operation of the type Real to each component of Left and the matching component of Right. The index range of the result is Left'Range. Constraint_Error is raised if Left'Length is not equal to Right'Length. 38/2

```
function "*" (Left, Right : Real_Vector) return Real'Base;
```

This operation returns the inner product of Left and Right. Constraint_Error is raised if Left'Length is not equal to Right'Length. This operation involves an inner product. 40/2

```
function "abs" (Right : Real_Vector) return Real'Base;
```

This operation returns the L2-norm of Right (the square root of the inner product of the vector with itself). 42/2

```
function "*" (Left : Real'Base; Right : Real_Vector) return Real_Vector;
```

This operation returns the result of multiplying each component of Right by the scalar Left using the "*" operation of the type Real. The index range of the result is Right'Range. 44/2

```
function "*" (Left : Real_Vector; Right : Real'Base) return Real_Vector;
function "/" (Left : Real_Vector; Right : Real'Base) return Real_Vector;
```

Each operation returns the result of applying the corresponding operation of the type Real to each component of Left and to the scalar Right. The index range of the result is Left'Range. 46/2

```
function Unit_Vector (Index : Integer;
                      Order : Positive;
                      First : Integer := 1) return Real_Vector;
```

This function returns a *unit vector* with Order components and a lower bound of First. All components are set to 0.0 except for the Index component which is set to 1.0. Constraint_Error is raised if Index < First, Index > First + Order - 1 or if First + Order - 1 > Integer'Last. 48/2

```
function "+" (Right : Real_Matrix) return Real_Matrix;
function "-" (Right : Real_Matrix) return Real_Matrix;
function "abs" (Right : Real_Matrix) return Real_Matrix;
```

Each operation returns the result of applying the corresponding operation of the type Real to each component of Right. The index ranges of the result are those of Right. 50/2

```

51/2   function Transpose (X : Real_Matrix) return Real_Matrix;
52/2   This function returns the transpose of a matrix X. The first and second index ranges of the result
      are X'Range(2) and X'Range(1) respectively.
53/2   function "+" (Left, Right : Real_Matrix) return Real_Matrix;
      function "-" (Left, Right : Real_Matrix) return Real_Matrix;
54/2   Each operation returns the result of applying the corresponding operation of the type Real to
      each component of Left and the matching component of Right. The index ranges of the result are
      those of Left. Constraint_Error is raised if Left'Length(1) is not equal to Right'Length(1) or
      Left'Length(2) is not equal to Right'Length(2).
55/2   function "*" (Left, Right : Real_Matrix) return Real_Matrix;
56/2   This operation provides the standard mathematical operation for matrix multiplication. The first
      and second index ranges of the result are Left'Range(1) and Right'Range(2) respectively.
      Constraint_Error is raised if Left'Length(2) is not equal to Right'Length(1). This operation
      involves inner products.
57/2   function "*" (Left, Right : Real_Vector) return Real_Matrix;
58/2   This operation returns the outer product of a (column) vector Left by a (row) vector Right using
      the operation "*" of the type Real for computing the individual components. The first and second
      index ranges of the result are Left'Range and Right'Range respectively.
59/2   function "*" (Left : Real_Vector; Right : Real_Matrix) return Real_Vector;
60/2   This operation provides the standard mathematical operation for multiplication of a (row) vector
      Left by a matrix Right. The index range of the (row) vector result is Right'Range(2).
      Constraint_Error is raised if Left'Length is not equal to Right'Length(1). This operation involves
      inner products.
61/2   function "*" (Left : Real_Matrix; Right : Real_Vector) return Real_Vector;
62/2   This operation provides the standard mathematical operation for multiplication of a matrix Left
      by a (column) vector Right. The index range of the (column) vector result is Left'Range(1).
      Constraint_Error is raised if Left'Length(2) is not equal to Right'Length. This operation involves
      inner products.
63/2   function "*" (Left : Real'Base; Right : Real_Matrix) return Real_Matrix;
64/2   This operation returns the result of multiplying each component of Right by the scalar Left using
      the "*" operation of the type Real. The index ranges of the result are those of Right.
65/2   function "*" (Left : Real_Matrix; Right : Real'Base) return Real_Matrix;
      function "/" (Left : Real_Matrix; Right : Real'Base) return Real_Matrix;
66/2   Each operation returns the result of applying the corresponding operation of the type Real to
      each component of Left and to the scalar Right. The index ranges of the result are those of Left.
67/2   function Solve (A : Real_Matrix; X : Real_Vector) return Real_Vector;
68/2   This function returns a vector Y such that X is (nearly) equal to A * Y. This is the standard
      mathematical operation for solving a single set of linear equations. The index range of the result
      is A'Range(2). Constraint_Error is raised if A'Length(1), A'Length(2), and X'Length are not
      equal. Constraint_Error is raised if the matrix A is ill-conditioned.
69/2   function Solve (A, X : Real_Matrix) return Real_Matrix;
70/2   This function returns a matrix Y such that X is (nearly) equal to A * Y. This is the standard
      mathematical operation for solving several sets of linear equations. The index ranges of the

```

result are A'Range(2) and X'Range(2). Constraint_Error is raised if A'Length(1), A'Length(2), and X'Length(1) are not equal. Constraint_Error is raised if the matrix A is ill-conditioned.

function Inverse (A : Real_Matrix) **return** Real_Matrix;

71/2

This function returns a matrix B such that A * B is (nearly) equal to the unit matrix. The index ranges of the result are A'Range(2) and A'Range(1). Constraint_Error is raised if A'Length(1) is not equal to A'Length(2). Constraint_Error is raised if the matrix A is ill-conditioned.

72/2

function Determinant (A : Real_Matrix) **return** Real'Base;

73/2

This function returns the determinant of the matrix A. Constraint_Error is raised if A'Length(1) is not equal to A'Length(2).

74/2

function Eigenvalues(A : Real_Matrix) **return** Real_Vector;

75/2

This function returns the eigenvalues of the symmetric matrix A as a vector sorted into order with the largest first. Constraint_Error is raised if A'Length(1) is not equal to A'Length(2). The index range of the result is A'Range(1). Argument_Error is raised if the matrix A is not symmetric.

76/2

procedure Eigensystem(A : in Real_Matrix;
Values : out Real_Vector;
Vectors : out Real_Matrix);

77/2

This procedure computes both the eigenvalues and eigenvectors of the symmetric matrix A. The out parameter Values is the same as that obtained by calling the function Eigenvalues. The out parameter Vectors is a matrix whose columns are the eigenvectors of the matrix A. The order of the columns corresponds to the order of the eigenvalues. The eigenvectors are normalized and mutually orthogonal (they are orthonormal), including when there are repeated eigenvalues. Constraint_Error is raised if A'Length(1) is not equal to A'Length(2). The index ranges of the parameter Vectors are those of A. Argument_Error is raised if the matrix A is not symmetric.

78/2

function Unit_Matrix (Order : Positive;
First_1, First_2 : Integer := 1) **return** Real_Matrix;

79/2

This function returns a square *unit matrix* with Order**2 components and lower bounds of First_1 and First_2 (for the first and second index ranges respectively). All components are set to 0.0 except for the main diagonal, whose components are set to 1.0. Constraint_Error is raised if First_1 + Order - 1 > Integer'Last or First_2 + Order - 1 > Integer'Last.

80/2

Implementation Requirements

Accuracy requirements for the subprograms Solve, Inverse, Determinant, Eigenvalues and Eigensystem are implementation defined.

81/2

For operations not involving an inner product, the accuracy requirements are those of the corresponding operations of the type Real in both the strict mode and the relaxed mode (see G.2).

82/2

For operations involving an inner product, no requirements are specified in the relaxed mode. In the strict mode the modulus of the absolute error of the inner product X^*Y shall not exceed $g * \text{abs}(X) * \text{abs}(Y)$ where g is defined as

83/2

$$g = X'Length * \text{Real}'\text{Machine_Radix}^{**}(1 - \text{Real}'\text{Model_Mantissa})$$

84/2

For the L2-norm, no accuracy requirements are specified in the relaxed mode. In the strict mode the relative error on the norm shall not exceed $g / 2.0 + 3.0 * \text{Real}'\text{Model_Epsilon}$ where g is defined as above.

85/2

Documentation Requirements

- 86/2 Implementations shall document any techniques used to reduce cancellation errors such as extended precision arithmetic.

Implementation Permissions

- 87/2 The nongeneric equivalent packages may, but need not, be actual instantiations of the generic package for the appropriate predefined type.

Implementation Advice

- 88/2 Implementations should implement the Solve and Inverse functions using established techniques such as LU decomposition with row interchanges followed by back and forward substitution. Implementations are recommended to refine the result by performing an iteration on the residuals; if this is done then it should be documented.
- 89/2 It is not the intention that any special provision should be made to determine whether a matrix is ill-conditioned or not. The naturally occurring overflow (including division by zero) which will result from executing these functions with an ill-conditioned matrix and thus raise Constraint_Error is sufficient.
- 90/2 The test that a matrix is symmetric should be performed by using the equality operator to compare the relevant components.

G.3.2 Complex Vectors and Matrices

Static Semantics

- 1/2 The generic library package Numerics.Generic_Complex_Arrays has the following declaration:

```

2/2   with Ada.Numerics.Generic_Real_Arrays, Ada.Numerics.Generic_Complex_Types;
generic
  with package Real_Arrays  is new
    Ada.Numerics.Generic_Real_Arrays  (<>);
  use Real_Arrays;
  with package Complex_Types is new
    Ada.Numerics.Generic_Complex_Types (Real);
  use Complex_Types;
package Ada.Numerics.Generic_Complex_Arrays is
  pragma Pure(Generic_Complex_Arrays);

3/2      -- Types
4/2      type Complex_Vector is array (Integer range <>) of Complex;
        type Complex_Matrix is array (Integer range <>,
                                      Integer range <>) of Complex;

5/2      -- Subprograms for Complex_Vector types
6/2      -- Complex_Vector selection, conversion and composition operations
7/2      function Re (X : Complex_Vector) return Real_Vector;
        function Im (X : Complex_Vector) return Real_Vector;
8/2      procedure Set_Re (X : in out Complex_Vector;
                           Re : in      Real_Vector);
      procedure Set_Im (X : in out Complex_Vector;
                           Im : in      Real_Vector);
9/2      function Compose_From_Cartesian (Re      : Real_Vector)
            return Complex_Vector;
      function Compose_From_Cartesian (Re, Im : Real_Vector)
            return Complex_Vector;

```

```

function Modulus  (X      : Complex_Vector) return Real_Vector;          10/2
function "abs"    (Right   : Complex_Vector) return Real_Vector;
                           renames Modulus;
function Argument (X      : Complex_Vector) return Real_Vector;
function Argument (X      : Complex_Vector;
                           Cycle   : Real'Base)      return Real_Vector;
function Compose_From_Polar (Modulus, Argument : Real_Vector)
                           return Complex_Vector;           11/2
function Compose_From_Polar (Modulus, Argument : Real_Vector;
                           Cycle       : Real'Base)
                           return Complex_Vector;
-- Complex_Vector arithmetic operations                                12/2
function "+"        (Right   : Complex_Vector) return Complex_Vector;
function "-"        (Right   : Complex_Vector) return Complex_Vector;
function Conjugate (X      : Complex_Vector) return Complex_Vector;
function "+"        (Left,   Right : Complex_Vector) return Complex_Vector; 14/2
function "-"        (Left,   Right : Complex_Vector) return Complex_Vector;
function "*"        (Left,   Right : Complex_Vector) return Complex;        15/2
function "abs"      (Right   : Complex_Vector) return Complex;            16/2
-- Mixed Real_Vector and Complex_Vector arithmetic operations             17/2
function "+"        (Left   : Real_Vector;
                           Right  : Complex_Vector) return Complex_Vector;
function "+"        (Left   : Complex_Vector;
                           Right  : Real_Vector)  return Complex_Vector;
function "-"        (Left   : Real_Vector;
                           Right  : Complex_Vector) return Complex_Vector;
function "-"        (Left   : Complex_Vector;
                           Right  : Real_Vector)  return Complex_Vector;
function "*"        (Left   : Real_Vector;     Right : Complex_Vector)
                           return Complex;                19/2
function "*"        (Left   : Complex_Vector; Right : Real_Vector)
                           return Complex;
-- Complex_Vector scaling operations                                     20/2
function "*"        (Left   : Complex;
                           Right  : Complex_Vector) return Complex_Vector;
function "*"        (Left   : Complex_Vector;
                           Right  : Complex)       return Complex_Vector;
function "/"        (Left   : Complex_Vector;
                           Right  : Complex)       return Complex_Vector;
function "*"        (Left   : Real'Base;
                           Right  : Complex_Vector) return Complex_Vector; 22/2
function "*"        (Left   : Complex_Vector;
                           Right  : Real'Base)      return Complex_Vector;
function "/"        (Left   : Complex_Vector;
                           Right  : Real'Base)      return Complex_Vector;
-- Other Complex_Vector operations                                    23/2
function Unit_Vector (Index : Integer;
                           Order : Positive;
                           First  : Integer := 1) return Complex_Vector;
-- Subprograms for Complex_Matrix types                               25/2
-- Complex_Matrix selection, conversion and composition operations    26/2
function Re (X : Complex_Matrix) return Real_Matrix;                  27/2
function Im (X : Complex_Matrix) return Real_Matrix;
procedure Set_Re (X : in out Complex_Matrix;
                           Re : in      Real_Matrix);
procedure Set_Im (X : in out Complex_Matrix;
                           Im : in      Real_Matrix);

```

```

29/2      function Compose_From_Cartesian (Re      : Real_Matrix)
           return Complex_Matrix;
      function Compose_From_Cartesian (Re, Im : Real_Matrix)
           return Complex_Matrix;
30/2      function Modulus   (X      : Complex_Matrix) return Real_Matrix;
      function "abs"     (Right : Complex_Matrix) return Real_Matrix
                           renames Modulus;
31/2      function Argument  (X      : Complex_Matrix) return Real_Matrix;
      function Argument  (X      : Complex_Matrix;
                           Cycle : Real'Base)      return Real_Matrix;
32/2      function Compose_From_Polar (Modulus, Argument : Real_Matrix)
           return Complex_Matrix;
      function Compose_From_Polar (Modulus, Argument : Real_Matrix;
                                   Cycle             : Real'Base)
           return Complex_Matrix;
33/2      -- Complex_Matrix arithmetic operations
34/2      function "+"        (Right : Complex_Matrix) return Complex_Matrix;
      function "-"        (Right : Complex_Matrix) return Complex_Matrix;
      function Conjugate (X      : Complex_Matrix) return Complex_Matrix;
      function Transpose (X      : Complex_Matrix) return Complex_Matrix;
35/2      function "+"        (Left, Right : Complex_Matrix) return Complex_Matrix;
      function "-"        (Left, Right : Complex_Matrix) return Complex_Matrix;
      function "*"        (Left, Right : Complex_Matrix) return Complex_Matrix;
36/2      function "*"        (Left, Right : Complex_Vector) return Complex_Matrix;
37/2      function "*"        (Left   : Complex_Vector;
                           Right  : Complex_Matrix) return Complex_Vector;
      function "*"        (Left   : Complex_Matrix;
                           Right  : Complex_Vector) return Complex_Vector;
38/2      -- Mixed Real_Matrix and Complex_Matrix arithmetic operations
39/2      function "+"        (Left   : Real_Matrix;
                           Right  : Complex_Matrix) return Complex_Matrix;
      function "+"        (Left   : Complex_Matrix;
                           Right  : Real_Matrix) return Complex_Matrix;
      function "-"        (Left   : Real_Matrix;
                           Right  : Complex_Matrix) return Complex_Matrix;
      function "-"        (Left   : Complex_Matrix;
                           Right  : Real_Matrix) return Complex_Matrix;
      function "*"        (Left   : Real_Matrix;
                           Right  : Complex_Matrix) return Complex_Matrix;
      function "*"        (Left   : Complex_Matrix;
                           Right  : Real_Matrix) return Complex_Matrix;
40/2      function "*"        (Left   : Real_Vector;
                           Right  : Complex_Vector) return Complex_Matrix;
      function "*"        (Left   : Complex_Vector;
                           Right  : Real_Vector) return Complex_Matrix;
41/2      function "*"        (Left   : Real_Vector;
                           Right  : Complex_Matrix) return Complex_Vector;
      function "*"        (Left   : Complex_Vector;
                           Right  : Real_Matrix) return Complex_Vector;
      function "*"        (Left   : Real_Matrix;
                           Right  : Complex_Vector) return Complex_Vector;
      function "*"        (Left   : Complex_Matrix;
                           Right  : Real_Vector) return Complex_Vector;
42/2      -- Complex_Matrix scaling operations
43/2      function "*"        (Left   : Complex;
                           Right  : Complex_Matrix) return Complex_Matrix;
      function "*"        (Left   : Complex_Matrix;
                           Right  : Complex)      return Complex_Matrix;
      function "/"        (Left   : Complex_Matrix;
                           Right  : Complex)      return Complex_Matrix;

```

```

function "*" (Left   : Real'Base;
              Right  : Complex_Matrix) return Complex_Matrix; 44/2
function "*" (Left   : Complex_Matrix;
              Right  : Real'Base)    return Complex_Matrix;
function "/" (Left   : Complex_Matrix;
              Right  : Real'Base)    return Complex_Matrix;

-- Complex_Matrix inversion and related operations 45/2
function Solve (A : Complex_Matrix; X : Complex_Vector)
                return Complex_Vector; 46/2
function Solve (A, X : Complex_Matrix) return Complex_Matrix;
function Inverse (A : Complex_Matrix) return Complex_Matrix;
function Determinant (A : Complex_Matrix) return Complex;

-- Eigenvalues and vectors of a Hermitian matrix 47/2
function Eigenvalues(A : Complex_Matrix) return Real_Vector; 48/2
procedure Eigensystem(A          : in Complex_Matrix;
                      Values     : out Real_Vector;
                      Vectors   : out Complex_Matrix); 49/2

-- Other Complex_Matrix operations 50/2
function Unit_Matrix (Order      : Positive;
                      First_1, First_2 : Integer := 1)
                      return Complex_Matrix; 51/2
end Ada.Numerics.Generic_Complex_Arrays; 52/2

```

The library package Numerics.Complex_Arrays is declared pure and defines the same types and subprograms as Numerics.Generic_Complex_Arrays, except that the predefined type Float is systematically substituted for Real'Base, and the Real_Vector and Real_Matrix types exported by Numerics.Real_Arrays are systematically substituted for Real_Vector and Real_Matrix, and the Complex type exported by Numerics.Complex_Types is systematically substituted for Complex, throughout. Nongeneric equivalents for each of the other predefined floating point types are defined similarly, with the names Numerics.Short_Complex_Arrays, Numerics.Long_Complex_Arrays, etc.

Two types are defined and exported by Numerics.Generic_Complex_Arrays. The composite type Complex_Vector is provided to represent a vector with components of type Complex; it is defined as an unconstrained one-dimensional array with an index of type Integer. The composite type Complex_Matrix is provided to represent a matrix with components of type Complex; it is defined as an unconstrained, two-dimensional array with indices of type Integer.

The effect of the various subprograms is as described below. In many cases they are described in terms of corresponding scalar operations in Numerics.Generic_Complex_Types. Any exception raised by those operations is propagated by the array subprogram. Moreover, any constraints on the parameters and the accuracy of the result for each individual component are as defined for the scalar operation.

In the case of those operations which are defined to *involve an inner product*, Constraint_Error may be raised if an intermediate result has a component outside the range of Real'Base even though the final mathematical result would not.

```

function Re (X : Complex_Vector) return Real_Vector;
function Im (X : Complex_Vector) return Real_Vector; 57/2

```

Each function returns a vector of the specified Cartesian components of X. The index range of the result is X'Range.

```

procedure Set_Re (X  : in out Complex_Vector; Re : in Real_Vector);
procedure Set_Im (X  : in out Complex_Vector; Im : in Real_Vector); 59/2

```

Each procedure replaces the specified (Cartesian) component of each of the components of X by the value of the matching component of Re or Im; the other (Cartesian) component of each of

the components is unchanged. Constraint_Error is raised if X'Length is not equal to Re'Length or Im'Length.

```
61/2  function Compose_From_Cartesian (Re      : Real_Vector)
      return Complex_Vector;
function Compose_From_Cartesian (Re, Im : Real_Vector)
      return Complex_Vector;
```

62/2 Each function constructs a vector of Complex results (in Cartesian representation) formed from given vectors of Cartesian components; when only the real components are given, imaginary components of zero are assumed. The index range of the result is Re'Range. Constraint_Error is raised if Re'Length is not equal to Im'Length.

```
63/2  function Modulus  (X      : Complex_Vector) return Real_Vector;
function "abs"    (Right : Complex_Vector) return Real_Vector
               renames Modulus;
function Argument (X      : Complex_Vector) return Real_Vector;
function Argument (X      : Complex_Vector,
                   Cycle   : Real'Base)      return Real_Vector;
```

64/2 Each function calculates and returns a vector of the specified polar components of X or Right using the corresponding function in numerics.generic_complex_types. The index range of the result is X'Range or Right'Range.

```
65/2  function Compose_From_Polar (Modulus, Argument : Real_Vector)
      return Complex_Vector;
function Compose_From_Polar (Modulus, Argument : Real_Vector;
                             Cycle           : Real'Base)
      return Complex_Vector;
```

66/2 Each function constructs a vector of Complex results (in Cartesian representation) formed from given vectors of polar components using the corresponding function in numerics.-generic_complex_types on matching components of Modulus and Argument. The index range of the result is Modulus'Range. Constraint_Error is raised if Modulus'Length is not equal to Argument'Length.

```
67/2  function "+" (Right : Complex_Vector) return Complex_Vector;
function "-" (Right : Complex_Vector) return Complex_Vector;
```

68/2 Each operation returns the result of applying the corresponding operation in numerics.-generic_complex_types to each component of Right. The index range of the result is Right'Range.

```
69/2  function Conjugate (X : Complex_Vector) return Complex_Vector;
```

70/2 This function returns the result of applying the appropriate function Conjugate in numerics.-generic_complex_types to each component of X. The index range of the result is X'Range.

```
71/2  function "+" (Left, Right : Complex_Vector) return Complex_Vector;
function "-" (Left, Right : Complex_Vector) return Complex_Vector;
```

72/2 Each operation returns the result of applying the corresponding operation in numerics.-generic_complex_types to each component of Left and the matching component of Right. The index range of the result is Left'Range. Constraint_Error is raised if Left'Length is not equal to Right'Length.

```
73/2  function "*" (Left, Right : Complex_Vector) return Complex;
```

74/2 This operation returns the inner product of Left and Right. Constraint_Error is raised if Left'Length is not equal to Right'Length. This operation involves an inner product.

function "abs" (Right : Complex_Vector) return Complex;	75/2
This operation returns the Hermitian L2-norm of Right (the square root of the inner product of the vector with its conjugate).	76/2
function "+" (Left : Real_Vector; Right : Complex_Vector) return Complex_Vector;	77/2
function "+" (Left : Complex_Vector; Right : Real_Vector) return Complex_Vector;	
function "-" (Left : Real_Vector; Right : Complex_Vector) return Complex_Vector;	
function "-" (Left : Complex_Vector; Right : Real_Vector) return Complex_Vector;	
Each operation returns the result of applying the corresponding operation in numerics.generic_complex_types to each component of Left and the matching component of Right. The index range of the result is Left'Range. Constraint_Error is raised if Left'Length is not equal to Right'Length.	78/2
function "*" (Left : Real_Vector; Right : Complex_Vector) return Complex;	79/2
function "*" (Left : Complex_Vector; Right : Real_Vector) return Complex;	
Each operation returns the inner product of Left and Right. Constraint_Error is raised if Left'Length is not equal to Right'Length. These operations involve an inner product.	80/2
function "*" (Left : Complex; Right : Complex_Vector) return Complex_Vector;	81/2
This operation returns the result of multiplying each component of Right by the complex number Left using the appropriate operation "*" in numerics.generic_complex_types. The index range of the result is Right'Range.	82/2
function "*" (Left : Complex_Vector; Right : Complex) return Complex_Vector;	83/2
function "/" (Left : Complex_Vector; Right : Complex) return Complex_Vector;	
Each operation returns the result of applying the corresponding operation in numerics.generic_complex_types to each component of the vector Left and the complex number Right. The index range of the result is Left'Range.	84/2
function "*" (Left : Real'Base; Right : Complex_Vector) return Complex_Vector;	85/2
This operation returns the result of multiplying each component of Right by the real number Left using the appropriate operation "*" in numerics.generic_complex_types. The index range of the result is Right'Range.	86/2
function "*" (Left : Complex_Vector; Right : Real'Base) return Complex_Vector;	87/2
function "/" (Left : Complex_Vector; Right : Real'Base) return Complex_Vector;	
Each operation returns the result of applying the corresponding operation in numerics.generic_complex_types to each component of the vector Left and the real number Right. The index range of the result is Left'Range.	88/2
function Unit_Vector (Index : Integer; Order : Positive; First : Integer := 1) return Complex_Vector;	89/2
This function returns a <i>unit vector</i> with Order components and a lower bound of First. All components are set to (0.0, 0.0) except for the Index component which is set to (1.0, 0.0). Constraint_Error is raised if Index < First, Index > First + Order - 1, or if First + Order - 1 > Integer'Last.	90/2

```

91/2   function Re (X : Complex_Matrix) return Real_Matrix;
function Im (X : Complex_Matrix) return Real_Matrix;

92/2   Each function returns a matrix of the specified Cartesian components of X. The index ranges of the result are those of X.

93/2   procedure Set_Re (X : in out Complex_Matrix; Re : in Real_Matrix);
procedure Set_Im (X : in out Complex_Matrix; Im : in Real_Matrix);

94/2   Each procedure replaces the specified (Cartesian) component of each of the components of X by the value of the matching component of Re or Im; the other (Cartesian) component of each of the components is unchanged. Constraint_Error is raised if X'Length(1) is not equal to Re'Length(1) or Im'Length(1) or if X'Length(2) is not equal to Re'Length(2) or Im'Length(2).

95/2   function Compose_From_Cartesian (Re      : Real_Matrix)
      return Complex_Matrix;
function Compose_From_Cartesian (Re, Im : Real_Matrix)
      return Complex_Matrix;

96/2   Each function constructs a matrix of Complex results (in Cartesian representation) formed from given matrices of Cartesian components; when only the real components are given, imaginary components of zero are assumed. The index ranges of the result are those of Re. Constraint_Error is raised if Re'Length(1) is not equal to Im'Length(1) or Re'Length(2) is not equal to Im'Length(2).

97/2   function Modulus  (X      : Complex_Matrix) return Real_Matrix;
function "abs"     (Right : Complex_Matrix) return Real_Matrix
                           renames Modulus;
function Argument (X      : Complex_Matrix) return Real_Matrix;
function Argument (X      : Complex_Matrix;
                   Cycle  : Real'Base)      return Real_Matrix;

98/2   Each function calculates and returns a matrix of the specified polar components of X or Right using the corresponding function in numerics.generic_complex_types. The index ranges of the result are those of X or Right.

99/2   function Compose_From_Polar (Modulus, Argument : Real_Matrix)
      return Complex_Matrix;
function Compose_From_Polar (Modulus, Argument : Real_Matrix;
                           Cycle           : Real'Base)
      return Complex_Matrix;

100/2  Each function constructs a matrix of Complex results (in Cartesian representation) formed from given matrices of polar components using the corresponding function in numerics.generic_complex_types on matching components of Modulus and Argument. The index ranges of the result are those of Modulus. Constraint_Error is raised if Modulus'Length(1) is not equal to Argument'Length(1) or Modulus'Length(2) is not equal to Argument'Length(2).

101/2  function "+"  (Right : Complex_Matrix) return Complex_Matrix;
function "-"  (Right : Complex_Matrix) return Complex_Matrix;

102/2  Each operation returns the result of applying the corresponding operation in numerics.generic_complex_types to each component of Right. The index ranges of the result are those of Right.

103/2  function Conjugate (X : Complex_Matrix) return Complex_Matrix;

104/2  This function returns the result of applying the appropriate function Conjugate in numerics.generic_complex_types to each component of X. The index ranges of the result are those of X.

```

```

function Transpose (X : Complex_Matrix) return Complex_Matrix; 105/2
  This function returns the transpose of a matrix X. The first and second index ranges of the result 106/2
  are X'Range(2) and X'Range(1) respectively.

function "+" (Left, Right : Complex_Matrix) return Complex_Matrix; 107/2
function "-" (Left, Right : Complex_Matrix) return Complex_Matrix; 107/2
  Each operation returns the result of applying the corresponding operation in numerics.- 108/2
  generic_complex_types to each component of Left and the matching component of Right. The 108/2
  index ranges of the result are those of Left. Constraint_Error is raised if Left'Length(1) is not 108/2
  equal to Right'Length(1) or Left'Length(2) is not equal to Right'Length(2).

function "*" (Left, Right : Complex_Matrix) return Complex_Matrix; 109/2
  This operation provides the standard mathematical operation for matrix multiplication. The first 110/2
  and second index ranges of the result are Left'Range(1) and Right'Range(2) respectively. 110/2
  Constraint_Error is raised if Left'Length(2) is not equal to Right'Length(1). This operation 110/2
  involves inner products.

function "*" (Left, Right : Complex_Vector) return Complex_Matrix; 111/2
  This operation returns the outer product of a (column) vector Left by a (row) vector Right using 112/2
  the appropriate operation "*" in numerics.generic_complex_types for computing the individual 112/2
  components. The first and second index ranges of the result are Left'Range and Right'Range 112/2
  respectively.

function "*" (Left : Complex_Vector;
               Right : Complex_Matrix) return Complex_Vector; 113/2
  This operation provides the standard mathematical operation for multiplication of a (row) vector 114/2
  Left by a matrix Right. The index range of the (row) vector result is Right'Range(2). 114/2
  Constraint_Error is raised if Left'Length is not equal to Right'Length(1). This operation involves 114/2
  inner products.

function "*" (Left : Complex_Matrix;
               Right : Complex_Vector) return Complex_Vector; 115/2
  This operation provides the standard mathematical operation for multiplication of a matrix Left 116/2
  by a (column) vector Right. The index range of the (column) vector result is Left'Range(1). 116/2
  Constraint_Error is raised if Left'Length(2) is not equal to Right'Length. This operation involves 116/2
  inner products.

function "+" (Left : Real_Matrix;
               Right : Complex_Matrix) return Complex_Matrix; 117/2
function "+" (Left : Complex_Matrix;
               Right : Real_Matrix) return Complex_Matrix;
function "-" (Left : Real_Matrix;
               Right : Complex_Matrix) return Complex_Matrix;
function "-" (Left : Complex_Matrix;
               Right : Real_Matrix) return Complex_Matrix;
  Each operation returns the result of applying the corresponding operation in numerics.- 118/2
  generic_complex_types to each component of Left and the matching component of Right. The 118/2
  index ranges of the result are those of Left. Constraint_Error is raised if Left'Length(1) is not 118/2
  equal to Right'Length(1) or Left'Length(2) is not equal to Right'Length(2).

```

```

119/2   function "*" (Left  : Real_Matrix;
                      Right : Complex_Matrix) return Complex_Matrix;
function "*" (Left  : Complex_Matrix;
                      Right : Real_Matrix)      return Complex_Matrix;

120/2   Each operation provides the standard mathematical operation for matrix multiplication. The first
        and second index ranges of the result are Left'Range(1) and Right'Range(2) respectively.
        Constraint_Error is raised if Left'Length(2) is not equal to Right'Length(1). These operations
        involve inner products.

121/2   function "*" (Left  : Real_Vector;
                      Right : Complex_Vector) return Complex_Matrix;
function "*" (Left  : Complex_Vector;
                      Right : Real_Vector)      return Complex_Matrix;

122/2   Each operation returns the outer product of a (column) vector Left by a (row) vector Right using
        the appropriate operation "*" in numerics.generic_complex_types for computing the individual
        components. The first and second index ranges of the result are Left'Range and Right'Range
        respectively.

123/2   function "*" (Left  : Real_Vector;
                      Right : Complex_Matrix) return Complex_Vector;
function "*" (Left  : Complex_Vector;
                      Right : Real_Matrix)      return Complex_Vector;

124/2   Each operation provides the standard mathematical operation for multiplication of a (row) vector
        Left by a matrix Right. The index range of the (row) vector result is Right'Range(2).
        Constraint_Error is raised if Left'Length is not equal to Right'Length(1). These operations
        involve inner products.

125/2   function "*" (Left  : Real_Matrix;
                      Right : Complex_Vector) return Complex_Vector;
function "*" (Left  : Complex_Matrix;
                      Right : Real_Vector)      return Complex_Vector;

126/2   Each operation provides the standard mathematical operation for multiplication of a matrix Left
        by a (column) vector Right. The index range of the (column) vector result is Left'Range(1).
        Constraint_Error is raised if Left'Length(2) is not equal to Right'Length. These operations
        involve inner products.

127/2   function "*" (Left : Complex; Right : Complex_Matrix) return Complex_Matrix;

128/2   This operation returns the result of multiplying each component of Right by the complex number
        Left using the appropriate operation "*" in numerics.generic_complex_types. The index ranges
        of the result are those of Right.

129/2   function "*" (Left : Complex_Matrix; Right : Complex) return Complex_Matrix;
function "/" (Left : Complex_Matrix; Right : Complex) return Complex_Matrix;

130/2   Each operation returns the result of applying the corresponding operation in numerics-
        generic_complex_types to each component of the matrix Left and the complex number Right.
        The index ranges of the result are those of Left.

131/2   function "*" (Left : Real'Base;
                      Right : Complex_Matrix) return Complex_Matrix;

132/2   This operation returns the result of multiplying each component of Right by the real number Left
        using the appropriate operation "*" in numerics.generic_complex_types. The index ranges of the
        result are those of Right.

```

function "*" (Left : Complex_Matrix; Right : Real'Base) return Complex_Matrix;	133/2
function "/" (Left : Complex_Matrix; Right : Real'Base) return Complex_Matrix;	
Each operation returns the result of applying the corresponding operation in numerics.-generic_complex_types to each component of the matrix Left and the real number Right. The index ranges of the result are those of Left.	134/2
function Solve (A : Complex_Matrix; X : Complex_Vector) return Complex_Vector;	135/2
This function returns a vector Y such that X is (nearly) equal to A * Y. This is the standard mathematical operation for solving a single set of linear equations. The index range of the result is A'Range(2). Constraint_Error is raised if A'Length(1), A'Length(2), and X'Length are not equal. Constraint_Error is raised if the matrix A is ill-conditioned.	136/2
function Solve (A, X : Complex_Matrix) return Complex_Matrix;	137/2
This function returns a matrix Y such that X is (nearly) equal to A * Y. This is the standard mathematical operation for solving several sets of linear equations. The index ranges of the result are A'Range(2) and X'Range(2). Constraint_Error is raised if A'Length(1), A'Length(2), and X'Length(1) are not equal. Constraint_Error is raised if the matrix A is ill-conditioned.	138/2
function Inverse (A : Complex_Matrix) return Complex_Matrix;	139/2
This function returns a matrix B such that A * B is (nearly) equal to the unit matrix. The index ranges of the result are A'Range(2) and A'Range(1). Constraint_Error is raised if A'Length(1) is not equal to A'Length(2). Constraint_Error is raised if the matrix A is ill-conditioned.	140/2
function Determinant (A : Complex_Matrix) return Complex;	141/2
This function returns the determinant of the matrix A. Constraint_Error is raised if A'Length(1) is not equal to A'Length(2).	142/2
function Eigenvalues(A : Complex_Matrix) return Real_Vector;	143/2
This function returns the eigenvalues of the Hermitian matrix A as a vector sorted into order with the largest first. Constraint_Error is raised if A'Length(1) is not equal to A'Length(2). The index range of the result is A'Range(1). Argument_Error is raised if the matrix A is not Hermitian.	144/2
procedure Eigensystem(A : in Complex_Matrix; Values : out Real_Vector; Vectors : out Complex_Matrix);	145/2
This procedure computes both the eigenvalues and eigenvectors of the Hermitian matrix A. The out parameter Values is the same as that obtained by calling the function Eigenvalues. The out parameter Vectors is a matrix whose columns are the eigenvectors of the matrix A. The order of the columns corresponds to the order of the eigenvalues. The eigenvectors are mutually orthonormal, including when there are repeated eigenvalues. Constraint_Error is raised if A'Length(1) is not equal to A'Length(2). The index ranges of the parameter Vectors are those of A. Argument_Error is raised if the matrix A is not Hermitian.	146/2
function Unit_Matrix (Order : Positive; First_1, First_2 : Integer := 1) return Complex_Matrix;	147/2
This function returns a square <i>unit matrix</i> with Order**2 components and lower bounds of First_1 and First_2 (for the first and second index ranges respectively). All components are set to	148/2

(0.0, 0.0) except for the main diagonal, whose components are set to (1.0, 0.0). Constraint_Error is raised if First_1 + Order - 1 > Integer'Last or First_2 + Order - 1 > Integer'Last.

Implementation Requirements

- 149/2 Accuracy requirements for the subprograms Solve, Inverse, Determinant, Eigenvalues and Eigensystem are implementation defined.
- 150/2 For operations not involving an inner product, the accuracy requirements are those of the corresponding operations of the type Real'Base and Complex in both the strict mode and the relaxed mode (see G.2).
- 151/2 For operations involving an inner product, no requirements are specified in the relaxed mode. In the strict mode the modulus of the absolute error of the inner product X^*Y shall not exceed $g * \text{abs}(X) * \text{abs}(Y)$ where g is defined as
 - 152/2
$$g = X\text{Length} * \text{Real}'\text{Machine}_\text{Radix}^{**}(1 - \text{Real}'\text{Model}_\text{Mantissa})$$
 for mixed complex and real operands
 - 153/2
$$g = \sqrt(2.0) * X\text{Length} * \text{Real}'\text{Machine}_\text{Radix}^{**}(1 - \text{Real}'\text{Model}_\text{Mantissa})$$
 for two complex operands
- 154/2 For the L2-norm, no accuracy requirements are specified in the relaxed mode. In the strict mode the relative error on the norm shall not exceed $g / 2.0 + 3.0 * \text{Real}'\text{Model}_\text{Epsilon}$ where g has the definition appropriate for two complex operands.

Documentation Requirements

- 155/2 Implementations shall document any techniques used to reduce cancellation errors such as extended precision arithmetic.

Implementation Permissions

- 156/2 The nongeneric equivalent packages may, but need not, be actual instantiations of the generic package for the appropriate predefined type.
- 157/2 Although many operations are defined in terms of operations from numerics.generic_complex_types, they need not be implemented by calling those operations provided that the effect is the same.

Implementation Advice

- 158/2 Implementations should implement the Solve and Inverse functions using established techniques. Implementations are recommended to refine the result by performing an iteration on the residuals; if this is done then it should be documented.
- 159/2 It is not the intention that any special provision should be made to determine whether a matrix is ill-conditioned or not. The naturally occurring overflow (including division by zero) which will result from executing these functions with an ill-conditioned matrix and thus raise Constraint_Error is sufficient.
- 160/2 The test that a matrix is Hermitian should use the equality operator to compare the real components and negation followed by equality to compare the imaginary components (see G.2.1).
- 161/2 Implementations should not perform operations on mixed complex and real operands by first converting the real operand to complex. See G.1.1.

Annex H (normative) High Integrity Systems

This Annex addresses requirements for high integrity systems (including safety-critical systems and security-critical systems). It provides facilities and specifies documentation requirements that relate to several needs:

- Understanding program execution; 2
- Reviewing object code; 3
- Restricting language constructs whose usage might complicate the demonstration of program correctness 4

Execution understandability is supported by pragma Normalize_Scalars, and also by requirements for the implementation to document the effect of a program in the presence of a bounded error or where the language rules leave the effect unspecified. 4.1

The pragmas Reviewable and Restrictions relate to the other requirements addressed by this Annex. 5

NOTES

- 1 The Valid attribute (see 13.9.2) is also useful in addressing these needs, to avoid problems that could otherwise arise from scalars that have values outside their declared range constraints. 6

H.1 Pragma Normalize_Scalars

This pragma ensures that an otherwise uninitialized scalar object is set to a predictable value, but out of range if possible. 1

Syntax

The form of a **pragma** Normalize_Scalars is as follows: 2

pragma Normalize_Scalars; 3

Post-Compilation Rules

Pragma Normalize_Scalars is a configuration pragma. It applies to all compilation_units included in a partition. 4

Documentation Requirements

If a **pragma** Normalize_Scalars applies, the implementation shall document the implicit initial values for scalar subtypes, and shall identify each case in which such a value is used and is not an invalid representation. 5/2

Implementation Advice

Whenever possible, the implicit initial values for a scalar subtype should be an invalid representation (see 13.9.1). 6/2

NOTES

- 2 The initialization requirement applies to uninitialized scalar objects that are subcomponents of composite objects, to allocated objects, and to stand-alone objects. It also applies to scalar **out** parameters. Scalar subcomponents of composite **out** parameters are initialized to the corresponding part of the actual, by virtue of 6.4.1. 7

- 8 3 The initialization requirement does not apply to a scalar for which pragma Import has been specified, since initialization
of an imported object is performed solely by the foreign language environment (see B.1).
- 9 4 The use of pragma Normalize_Scalars in conjunction with Pragma Restrictions(No_Exceptions) may result in
erroneous execution (see H.4).

H.2 Documentation of Implementation Decisions

Documentation Requirements

- 1 The implementation shall document the range of effects for each situation that the language rules identify
as either a bounded error or as having an unspecified effect. If the implementation can constrain the effects
of erroneous execution for a given construct, then it shall document such constraints. The documentation
might be provided either independently of any compilation unit or partition, or as part of an annotated
listing for a given unit or partition. See also 1.1.3, and 1.1.2.

NOTES

- 2 5 Among the situations to be documented are the conventions chosen for parameter passing, the methods used for the
management of run-time storage, and the method used to evaluate numeric expressions if this involves extended range or
extra precision.

H.3 Reviewable Object Code

- 1 Object code review and validation are supported by pragmas Reviewable and Inspection_Point.

H.3.1 Pragma Reviewable

- 1 This pragma directs the implementation to provide information to facilitate analysis and review of a
program's object code, in particular to allow determination of execution time and storage usage and to
identify the correspondence between the source and object programs.

Syntax

- 2 The form of a **pragma** Reviewable is as follows:
- 3 **pragma** Reviewable;

Post-Compilation Rules

- 4 Pragma Reviewable is a configuration pragma. It applies to all compilation_units included in a partition.

Implementation Requirements

- 5 The implementation shall provide the following information for any compilation unit to which such a
pragma applies:
- 6 • Where compiler-generated run-time checks remain;
- 7 • An identification of any construct with a language-defined check that is recognized prior to run
time as certain to fail if executed (even if the generation of run-time checks has been
suppressed);
- 8/2 • For each read of a scalar object, an identification of the read as either “known to be initialized,”
or “possibly uninitialized,” independent of whether pragma Normalize_Scalars applies;
- 9 • Where run-time support routines are implicitly invoked;

- An object code listing, including:
 - Machine instructions, with relative offsets;
 - Where each data object is stored during its lifetime;
 - Correspondence with the source program, including an identification of the code produced per declaration and per statement.
- An identification of each construct for which the implementation detects the possibility of erroneous execution;
- For each subprogram, block, task, or other construct implemented by reserving and subsequently freeing an area on a run-time stack, an identification of the length of the fixed-size portion of the area and an indication of whether the non-fixed size portion is reserved on the stack or in a dynamically-managed storage region.

The implementation shall provide the following information for any partition to which the pragma applies:

- An object code listing of the entire partition, including initialization and finalization code as well as run-time system components, and with an identification of those instructions and data that will be relocated at load time;
- A description of the run-time model relevant to the partition.

The implementation shall provide control- and data-flow information, both within each compilation unit and across the compilation units of the partition.

Implementation Advice

The implementation should provide the above information in both a human-readable and machine-readable form, and should document the latter so as to ease further processing by automated tools.

Object code listings should be provided both in a symbolic format and also in an appropriate numeric format (such as hexadecimal or octal).

NOTES

6 The order of elaboration of library units will be documented even in the absence of `pragma Reviewable` (see 10.2).

H.3.2 Pragma Inspection_Point

An occurrence of a `pragma Inspection_Point` identifies a set of objects each of whose values is to be available at the point(s) during program execution corresponding to the position of the pragma in the compilation unit. The purpose of such a pragma is to facilitate code validation.

Syntax

The form of a `pragma Inspection_Point` is as follows:

`pragma Inspection_Point[(object_name {, object_name})];`

Legality Rules

A `pragma Inspection_Point` is allowed wherever a `declarative_item` or `statement` is allowed. Each *object_name* shall statically denote the declaration of an object.

Static Semantics

An *inspection point* is a point in the object code corresponding to the occurrence of a `pragma Inspection_Point` in the compilation unit. An object is *inspectable* at an inspection point if the corresponding pragma

Inspection_Point either has an argument denoting that object, or has no arguments and the declaration of the object is visible at the inspection point.

Dynamic Semantics

- 6 Execution of a pragma Inspection_Point has no effect.

Implementation Requirements

- 7 Reaching an inspection point is an external interaction with respect to the values of the inspectable objects at that point (see 1.1.3).

Documentation Requirements

- 8 For each inspection point, the implementation shall identify a mapping between each inspectable object and the machine resources (such as memory locations or registers) from which the object's value can be obtained.

NOTES

- 9/2 7 The implementation is not allowed to perform “dead store elimination” on the last assignment to a variable prior to a point where the variable is inspectable. Thus an inspection point has the effect of an implicit read of each of its inspectable objects.
- 10 8 Inspection points are useful in maintaining a correspondence between the state of the program in source code terms, and the machine state during the program's execution. Assertions about the values of program objects can be tested in machine terms at inspection points. Object code between inspection points can be processed by automated tools to verify programs mechanically.
- 11 9 The identification of the mapping from source program objects to machine resources is allowed to be in the form of an annotated object listing, in human-readable or tool-processable form.

H.4 High Integrity Restrictions

- 1 This clause defines restrictions that can be used with pragma Restrictions (see 13.12); these facilitate the demonstration of program correctness by allowing tailored versions of the run-time system.

Static Semantics

- 2/2 *This paragraph was deleted.*

- 3/2 The following *restriction_identifiers* are language defined:

4 **Tasking-related restriction:**

5 **No_Protected_Types**

There are no declarations of protected types or protected objects.

6 **Memory-management related restrictions:**

7 **No_Allocators**

There are no occurrences of an allocator.

8/1 **No_Local_Allocators**

Allocators are prohibited in subprograms, generic subprograms, tasks, and entry bodies.

- 9/2 *This paragraph was deleted.*

Immediate_Reclamation	10
Except for storage occupied by objects created by allocators and not deallocated via unchecked deallocation, any storage reserved at run time for an object is immediately reclaimed when the object no longer exists.	
Exception-related restriction:	11
No_Exceptions	12
Raise_statements and exception_handlers are not allowed. No language-defined run-time checks are generated; however, a run-time check performed automatically by the hardware is permitted.	
Other restrictions:	13
No_Floating_Point	14
Uses of predefined floating point types and operations, and declarations of new floating point types, are not allowed.	
No_Fixed_Point	15
Uses of predefined fixed point types and operations, and declarations of new fixed point types, are not allowed.	
<i>This paragraph was deleted.</i>	16/2
No_Access_Subprograms	17
The declaration of access-to-subprogram types is not allowed.	
No_Unchecked_Access	18
The Unchecked_Access attribute is not allowed.	
No_Dispatch	19
Occurrences of T'Class are not allowed, for any (tagged) subtype T.	
No_IO	20/2
Semantic dependence on any of the library units Sequential_IO, Direct_IO, Text_IO, Wide_Text_IO, Wide_Wide_Text_IO, or Stream_IO is not allowed.	
No_Delay	21
Delay_Statements and semantic dependence on package Calendar are not allowed.	
No_Recursion	22
As part of the execution of a subprogram, the same subprogram is not invoked.	
No_Reentrancy	23
During the execution of a subprogram by a task, no other task invokes the same subprogram.	

Implementation Requirements

An implementation of this Annex shall support:	23.1/2
• the restrictions defined in this subclause; and	23.2/2
• the following restrictions defined in D.7: No_Task_Hierarchy, No_Abort_Statement, No_Implicit_Heap_Allocation; and	23.3/2
• the pragma Profile(Ravenscar); and	23.4/2
• the following uses of <i>restriction_parameter_identifiers</i> defined in D.7, which are checked prior to program execution:	23.5/2
• Max_Task_Entries => 0,	23.6/2
• Max_Asynchronous_Select_Nesting => 0, and	23.7/2
• Max_Tasks => 0.	23.8/2

- 24 If an implementation supports **pragma** Restrictions for a particular argument, then except for the restrictions No_Unchecked_Deallocation, No_Unchecked_Conversion, No_Access_Subprograms, and No_Unchecked_Access, the associated restriction applies to the run-time system.

Documentation Requirements

- 25 If a **pragma** Restrictions(No_Exceptions) is specified, the implementation shall document the effects of all constructs where language-defined checks are still performed automatically (for example, an overflow check performed by the processor).

Erroneous Execution

- 26 Program execution is erroneous if **pragma** Restrictions(No_Exceptions) has been specified and the conditions arise under which a generated language-defined run-time check would fail.
- 27 Program execution is erroneous if **pragma** Restrictions(No_Recursion) has been specified and a subprogram is invoked as part of its own execution, or if **pragma** Restrictions(No_Reentrancy) has been specified and during the execution of a subprogram by a task, another task invokes the same subprogram.

NOTES

- 28/2 10 Uses of *restriction_parameter_identifier* No_Dependence defined in 13.12.1: No_Dependence => Ada.Unchecked_Deallocation and No_Dependence => Ada.Unchecked_Conversion may be appropriate for high-integrity systems. Other uses of No_Dependence can also be appropriate for high-integrity systems.

H.5 Pragma Detect_Blocking

- 1/2 The following **pragma** forces an implementation to detect potentially blocking operations within a protected operation.

Syntax

- 2/2 The form of a **pragma** Detect_Blocking is as follows:
- 3/2 **pragma** Detect_Blocking;

Post-Compilation Rules

- 4/2 A **pragma** Detect_Blocking is a configuration pragma.

Dynamic Semantics

- 5/2 An implementation is required to detect a potentially blocking operation within a protected operation, and to raise Program_Error (see 9.5.1).

Implementation Permissions

- 6/2 An implementation is allowed to reject a compilation_unit if a potentially blocking operation is present directly within an entry_body or the body of a protected subprogram.

NOTES

- 7/2 11 An operation that causes a task to be blocked within a foreign language domain is not defined to be potentially blocking, and need not be detected.

H.6 Pragma Partition_Elaboration_Policy

This clause defines a **pragma** for user control over elaboration policy.

1/2

Syntax

The form of a **pragma** Partition_Elaboration_Policy is as follows:

2/2

pragma Partition_Elaboration_Policy (*policy_identifier*);

3/2

The *policy_identifier* shall be either Sequential, Concurrent or an implementation-defined identifier.

4/2

Post-Compilation Rules

A **pragma** Partition_Elaboration_Policy is a configuration pragma. It specifies the elaboration policy for a partition. At most one elaboration policy shall be specified for a partition.

5/2

If the Sequential policy is specified for a partition then **pragma** Restrictions (No_Task_Hierarchy) shall also be specified for the partition.

6/2

Dynamic Semantics

Notwithstanding what this International Standard says elsewhere, this **pragma** allows partition elaboration rules concerning task activation and interrupt attachment to be changed. If the *policy_identifier* is Concurrent, or if there is no **pragma** Partition_Elaboration_Policy defined for the partition, then the rules defined elsewhere in this Standard apply.

7/2

If the partition elaboration policy is Sequential, then task activation and interrupt attachment are performed in the following sequence of steps:

8/2

- The activation of all library-level tasks and the attachment of interrupt handlers are deferred until all library units are elaborated.
- 9/2
- The interrupt handlers are attached by the environment task.
- 10/2
- The environment task is suspended while the library-level tasks are activated.
- 11/2
- The environment task executes the main subprogram (if any) concurrently with these executing tasks.
- 12/2

If several dynamic interrupt handler attachments for the same interrupt are deferred, then the most recent call of Attach_Handler or Exchange_Handler determines which handler is attached.

13/2

If any deferred task activation fails, Tasking_Error is raised at the beginning of the sequence of statements of the body of the environment task prior to calling the main subprogram.

14/2

Implementation Advice

If the partition elaboration policy is Sequential and the Environment task becomes permanently blocked during elaboration then the partition is deadlocked and it is recommended that the partition be immediately terminated.

15/2

Implementation Permissions

If the partition elaboration policy is Sequential and any task activation fails then an implementation may immediately terminate the active partition to mitigate the hazard posed by continuing to execute with a subset of the tasks being active.

16/2

NOTES

17/2

- 12 If any deferred task activation fails, the environment task is unable to handle the `Tasking_Error` exception and completes immediately. By contrast, if the partition elaboration policy is `Concurrent`, then this exception could be handled within a library unit.

Annex J (normative) Obsolescent Features

This Annex contains descriptions of features of the language whose functionality is largely redundant with other features defined by this International Standard. Use of these features is not recommended in newly written programs. Use of these features can be prevented by using pragma Restrictions (No_Obsolescent_Features), see 13.12.1.

J.1 Renamings of Ada 83 Library Units

Static Semantics

The following library_unit_renaming_declarations exist:

```

1   with Ada.Unchecked_Conversion;
2   generic function Unchecked_Conversion renames Ada.Unchecked_Conversion;
3   with Ada.Unchecked_Deallocation;
4   generic procedure Unchecked_Deallocation renames Ada.Unchecked_Deallocation;
5   with Ada.Sequential_IO;
6   generic package Sequential_IO renames Ada.Sequential_IO;
7   with Ada.Direct_IO;
8   generic package Direct_IO renames Ada.Direct_IO;
9   with Ada.Text_IO;
10  package Text_IO renames Ada.Text_IO;
11  with Ada.IO_Exceptions;
12  package IO_Exceptions renames Ada.IO_Exceptions;
13  with Ada.Calendar;
14  package Calendar renames Ada.Calendar;
15  with System.Machine_Code;
16  package Machine_Code renames System.Machine_Code; -- If supported.

```

Implementation Requirements

The implementation shall allow the user to replace these renamings.

J.2 Allowed Replacements of Characters

Syntax

The following replacements are allowed for the vertical line, number sign, and quotation mark characters:

- A vertical line character (|) can be replaced by an exclamation mark (!) where used as a delimiter.
- The number sign characters (#) of a based_literal can be replaced by colons (:) provided that the replacement is done for both occurrences.
- The quotation marks ("") used as string brackets at both ends of a string literal can be replaced by percent signs (%) provided that the enclosed sequence of characters contains no quotation mark, and provided that both string brackets are replaced. Any percent sign within the sequence of characters shall then be doubled and each such doubled percent sign is interpreted as a single percent sign character value.

- 5 These replacements do not change the meaning of the program.

J.3 Reduced Accuracy Subtypes

- 1 A `digits_constraint` may be used to define a floating point subtype with a new value for its requested decimal precision, as reflected by its `Digits` attribute. Similarly, a `delta_constraint` may be used to define an ordinary fixed point subtype with a new value for its `delta`, as reflected by its `Delta` attribute.

Syntax

- 2 `delta_constraint ::= delta static_expression [range_constraint]`

Name Resolution Rules

- 3 The expression of a `delta_constraint` is expected to be of any real type.

Legality Rules

- 4 The expression of a `delta_constraint` shall be static.
- 5 For a `subtype_indication` with a `delta_constraint`, the `subtype_mark` shall denote an ordinary fixed point subtype.
- 6 For a `subtype_indication` with a `digits_constraint`, the `subtype_mark` shall denote either a decimal fixed point subtype or a floating point subtype (notwithstanding the rule given in 3.5.9 that only allows a decimal fixed point subtype).

Static Semantics

- 7 A `subtype_indication` with a `subtype_mark` that denotes an ordinary fixed point subtype and a `delta_constraint` defines an ordinary fixed point subtype with a `delta` given by the value of the expression of the `delta_constraint`. If the `delta_constraint` includes a `range_constraint`, then the ordinary fixed point subtype is constrained by the `range_constraint`.
- 8 A `subtype_indication` with a `subtype_mark` that denotes a floating point subtype and a `digits_constraint` defines a floating point subtype with a requested decimal precision (as reflected by its `Digits` attribute) given by the value of the expression of the `digits_constraint`. If the `digits_constraint` includes a `range_constraint`, then the floating point subtype is constrained by the `range_constraint`.

Dynamic Semantics

- 9 A `delta_constraint` is *compatible* with an ordinary fixed point subtype if the value of the expression is no less than the `delta` of the subtype, and the `range_constraint`, if any, is compatible with the subtype.
- 10 A `digits_constraint` is *compatible* with a floating point subtype if the value of the expression is no greater than the requested decimal precision of the subtype, and the `range_constraint`, if any, is compatible with the subtype.
- 11 The elaboration of a `delta_constraint` consists of the elaboration of the `range_constraint`, if any.

J.4 The Constrained Attribute

Static Semantics

For every private subtype S, the following attribute is defined:

S'Constrained

Yields the value False if S denotes an unconstrained nonformal private subtype with discriminants; also yields the value False if S denotes a generic formal private subtype, and the associated actual subtype is either an unconstrained subtype with discriminants or an unconstrained array subtype; yields the value True otherwise. The value of this attribute is of the predefined subtype Boolean.

J.5 ASCII

Static Semantics

The following declaration exists in the declaration of package Standard:

```

package ASCII is
  -- Control characters:
  NUL   : constant Character := nul;      SOH   : constant Character := soh;
  STX   : constant Character := stx;      ETX   : constant Character := etx;
  EOT   : constant Character := eot;      ENQ   : constant Character := enq;
  ACK   : constant Character := ack;      BEL   : constant Character := bel;
  BS    : constant Character := bs;       HT    : constant Character := ht;
  LF    : constant Character := lf;       VT    : constant Character := vt;
  FF    : constant Character := ff;       CR    : constant Character := cr;
  SO    : constant Character := so;       SI    : constant Character := si;
  DLE   : constant Character := dle;     DC1   : constant Character := dc1;
  DC2   : constant Character := dc2;     DC3   : constant Character := dc3;
  DC4   : constant Character := dc4;     NAK   : constant Character := nak;
  SYN   : constant Character := syn;     ETB   : constant Character := etb;
  CAN   : constant Character := can;     EM    : constant Character := em;
  SUB   : constant Character := sub;     ESC   : constant Character := esc;
  FS    : constant Character := fs;      GS    : constant Character := gs;
  RS    : constant Character := rs;      US    : constant Character := us;
  DEL   : constant Character := del;

  -- Other characters:
  Exclam : constant Character:= '!'; Quotation : constant Character:= '"';
  Sharp  : constant Character:= '#'; Dollar  : constant Character:= '$';
  Percent : constant Character:= '%'; Ampersand : constant Character:= '&';
  Colon  : constant Character:= ':'; Semicolon : constant Character:= ';';
  Query  : constant Character:= '?'; At_Sign : constant Character:= '@';
  L_Bracket: constant Character:= '['; Back_Slash: constant Character:= '\';
  R_Bracket: constant Character:= ']'; Circumflex: constant Character:= '^';
  Underline: constant Character:= '_'; Grave   : constant Character:= ``;
  L_Brace : constant Character:= '{'; Bar     : constant Character:= '|';
  R_Brace : constant Character:= '}'; Tilde   : constant Character:= '~';

  -- Lower case letters:
  LC_A: constant Character:= 'a';
  ...
  LC_Z: constant Character:= 'z';

end ASCII;
```

J.6 Numeric_Error

Static Semantics

- 1 The following declaration exists in the declaration of package Standard:

```
Numeric_Error : exception renames Constraint_Error;
```

J.7 At Clauses

Syntax

- 1 `at_clause ::= for direct_name use at expression;`

Static Semantics

- 2 An `at_clause` of the form “`for x use at y;`” is equivalent to an `attribute_definition_clause` of the form “`for x'Address use y;`”.

J.7.1 Interrupt Entries

- 1 Implementations are permitted to allow the attachment of task entries to interrupts via the address clause. Such an entry is referred to as an *interrupt entry*.
- 2 The address of the task entry corresponds to a hardware interrupt in an implementation-defined manner. (See Ada.Interrupts.Reference in C.3.2.)

Static Semantics

- 3 The following attribute is defined:

- 4 For any task entry X:

- 5 `X'Address` For a task entry whose address is specified (an *interrupt entry*), the value refers to the corresponding hardware interrupt. For such an entry, as for any other task entry, the meaning of this value is implementation defined. The value of this attribute is of the type of the subtype System.Address.

- 6 Address may be specified for single entries via an `attribute_definition_clause`.

Dynamic Semantics

- 7 As part of the initialization of a task object, the address clause for an interrupt entry is elaborated, which evaluates the `expression` of the address clause. A check is made that the address specified is associated with some interrupt to which a task entry may be attached. If this check fails, `Program_Error` is raised. Otherwise, the interrupt entry is attached to the interrupt associated with the specified address.

- 8 Upon finalization of the task object, the interrupt entry, if any, is detached from the corresponding interrupt and the default treatment is restored.

- 9 While an interrupt entry is attached to an interrupt, the interrupt is reserved (see C.3).

- 10 An interrupt delivered to a task entry acts as a call to the entry issued by a hardware task whose priority is in the `System.Interrupt_Priority` range. It is implementation defined whether the call is performed as an ordinary entry call, a timed entry call, or a conditional entry call; which kind of call is performed can depend on the specific interrupt.

Bounded (Run-Time) Errors

It is a bounded error to evaluate E'Caller (see C.7.1) in an **accept_statement** for an interrupt entry. The possible effects are the same as for calling Current_Task from an entry body.

Documentation Requirements

The implementation shall document to which interrupts a task entry may be attached.

The implementation shall document whether the invocation of an interrupt entry has the effect of an ordinary entry call, conditional call, or a timed call, and whether the effect varies in the presence of pending interrupts.

Implementation Permissions

The support for this subclause is optional.

Interrupts to which the implementation allows a task entry to be attached may be designated as reserved for the entire duration of program execution; that is, not just when they have an interrupt entry attached to them.

Interrupt entry calls may be implemented by having the hardware execute directly the appropriate **accept_statement**. Alternatively, the implementation is allowed to provide an internal interrupt handler to simulate the effect of a normal task calling the entry.

The implementation is allowed to impose restrictions on the specifications and bodies of tasks that have interrupt entries.

It is implementation defined whether direct calls (from the program) to interrupt entries are allowed.

If a **select_statement** contains both a **terminate_alternative** and an **accept_alternative** for an interrupt entry, then an implementation is allowed to impose further requirements for the selection of the **terminate_alternative** in addition to those given in 9.3.

NOTES

1 Queued interrupts correspond to ordinary entry calls. Interrupts that are lost if not immediately processed correspond to conditional entry calls. It is a consequence of the priority rules that an **accept_statement** executed in response to an interrupt can be executed with the active priority at which the hardware generates the interrupt, taking precedence over lower priority tasks, without a scheduling action.

2 Control information that is supplied upon an interrupt can be passed to an associated interrupt entry as one or more parameters of mode **in**.

Examples

Example of an interrupt entry:

```
task Interrupt_Handler is
  entry Done;
    for Done'Address use
      Ada.Interrupts.Reference(Ada.Interrupts.Names.Device_Done);
end Interrupt_Handler;
```

J.8 Mod Clauses

Syntax

1 mod_clause ::= **at mod** static_expression;

Static Semantics

2 A record_representation_clause of the form:

3 **for** r **use**
 record at mod a
 ...
 end record;

4 is equivalent to:

5 **for** r'Alignment **use** a;
 for r **use**
 record
 ...
 end record;

J.9 The Storage_Size Attribute

Static Semantics

1 For any task subtype T, the following attribute is defined:

2 T'Storage_Size

Denotes an implementation-defined value of type *universal_integer* representing the number of storage elements reserved for a task of the subtype T.

3/2 Storage_Size may be specified for a task first subtype that is not an interface via an attribute_definition_clause.

J.10 Specific Suppression of Checks

1/2 Pragma Suppress can be used to suppress checks on specific entities.

Syntax

2/2 The form of a specific Suppress pragma is as follows:

3/2 **pragma** Suppress(identifier, [On =>] name);

Legality Rules

4/2 The identifier shall be the name of a check (see 11.5). The name shall statically denote some entity.

5/2 For a specific Suppress pragma that is immediately within a package_specification, the name shall denote an entity (or several overloaded subprograms) declared immediately within the package_specification.

Static Semantics

6/2 A specific Suppress pragma applies to the named check from the place of the pragma to the end of the innermost enclosing declarative region, or, if the pragma is given in a package_specification, to the end of the scope of the named entity. The pragma applies only to the named entity, or, for a subtype, on objects and values of its type. A specific Suppress pragma suppresses the named check for any entities to

which it applies (see 11.5). Which checks are associated with a specific entity is not defined by this International Standard.

Implementation Permissions

An implementation is allowed to place restrictions on specific Suppress pragmas.

7/2

NOTES

3 An implementation may support a similar On parameter on pragma Unsuppress (see 11.5).

8/2

J.11 The Class Attribute of Untagged Incomplete Types

Static Semantics

For the first subtype S of a type T declared by an incomplete_type_declaration that is not tagged, the following attribute is defined:

1/2

S'Class Denotes the first subtype of the incomplete class-wide type rooted at T. The completion of T shall declare a tagged type. Such an attribute reference shall occur in the same library unit as the incomplete_type_declaration.

2/2

J.12 Pragma Interface

Syntax

In addition to an identifier, the reserved word **interface** is allowed as a pragma name, to provide compatibility with a prior edition of this International Standard.

1/2

J.13 Dependence Restriction Identifiers

The following restrictions involve dependence on specific language-defined units. The more general restriction No_Dependence (see 13.12.1) should be used for this purpose.

1/2

Static Semantics

The following *restriction_identifiers* exist:

2/2

No_Asyncronous_Control

3/2

Semantic dependence on the predefined package Asynchronous_Task_Control is not allowed.

No_Unchecked_Conversion

4/2

Semantic dependence on the predefined generic function Unchecked_Conversion is not allowed.

No_Unchecked_Deallocation

5/2

Semantic dependence on the predefined generic procedure Unchecked_Deallocation is not allowed.

J.14 Character and Wide_Character Conversion Functions

Static Semantics

1/2 The following declarations exist in the declaration of package Ada.Characters.Handling:

```

2/2      function Is_Character (Item : in Wide_Character) return Boolean
             renames Conversions.Is_Character;
      function Is_String    (Item : in Wide_String)      return Boolean
             renames Conversions.Is_String;

3/2      function To_Character (Item      : in Wide_Character;
                           Substitute : in Character := ' ')
             return Character
             renames Conversions.To_Character;

4/2      function To_String     (Item      : in Wide_String;
                           Substitute : in Character := ' ')
             return String
             renames Conversions.To_String;

5/2      function To_Wide_Character (Item : in Character) return Wide_Character
             renames Conversions.To_Wide_Character;

6/2      function To_Wide_String   (Item : in String)      return Wide_String
             renames Conversions.To_Wide_String;
```

Annex K (informative) Language-Defined Attributes

This annex summarizes the definitions given elsewhere of the language-defined attributes.

P'Access	For a prefix P that denotes a subprogram:	1
	P'Access yields an access value that designates the subprogram denoted by P. The type of P'Access is an access-to-subprogram type (S), as determined by the expected type. See 3.10.2.	2
X'Access	For a prefix X that denotes an aliased view of an object:	3
	X'Access yields an access value that designates the object denoted by X. The type of X'Access is an access-to-object type, as determined by the expected type. The expected type shall be a general access type. See 3.10.2.	4
X'Address	For a prefix X that denotes an object, program unit, or label:	5
	Denotes the address of the first of the storage elements allocated to X. For a program unit or label, this value refers to the machine code associated with the corresponding body or statement. The value of this attribute is of type System.Address. See 13.3.	6/1
S'Adjacent	For every subtype S of a floating point type T:	7
	S'Adjacent denotes a function with the following specification:	8
	<code>function S'Adjacent (X, Towards : T) return T</code>	9
	If <i>Towards</i> = X, the function yields X; otherwise, it yields the machine number of the type T adjacent to X in the direction of <i>Towards</i> , if that machine number exists. If the result would be outside the base range of S, Constraint_Error is raised. When TSigned_Zeros is True, a zero result has the sign of X. When <i>Towards</i> is zero, its sign has no bearing on the result. See A.5.3.	10
S'Aft	For every fixed point subtype S:	11
	S'Aft yields the number of decimal digits needed after the decimal point to accommodate the <i>delta</i> of the subtype S, unless the <i>delta</i> of the subtype S is greater than 0.1, in which case the attribute yields the value one. (S'Aft is the smallest positive integer N for which $(10^{**N}) * S'\Delta$ is greater than or equal to one.) The value of this attribute is of the type universal_integer. See 3.5.10.	12
S'Alignment	For every subtype S:	13.1/2
	The value of this attribute is of type universal_integer, and nonnegative.	13.2/2
	For an object X of subtype S, if S'Alignment is not zero, then X'Alignment is a nonzero integral multiple of S'Alignment unless specified otherwise by a representation item. See 13.3.	13.3/2
X'Alignment	For a prefix X that denotes an object:	14/1
	The value of this attribute is of type universal_integer, and nonnegative; zero means that the object is not necessarily aligned on a storage element boundary. If X'Alignment is not zero, then X is aligned on a storage unit boundary and X'Address is an integral multiple of X'Alignment (that is, the Address modulo the Alignment is zero).	15
	See 13.3.	16/2

- 17 S'Base For every scalar subtype S:
 18 S'Base denotes an unconstrained subtype of the type of S. This unconstrained subtype is called the *base subtype* of the type. See 3.5.
- 19 S'Bit_Order For every specific record subtype S:
 20 Denotes the bit ordering for the type of S. The value of this attribute is of type System.Bit_Order. See 13.5.3.
- 21/1 P'Body_Version
 22 For a prefix P that statically denotes a program unit:
 23 Yields a value of the predefined type String that identifies the version of the compilation unit that contains the body (but not any subunits) of the program unit. See E.3.
- 24 T'Callable For a prefix T that is of a task type (after any implicit dereference):
 25 Yields the value True when the task denoted by T is *callable*, and False otherwise; See 9.9.
- 26 E'Caller For a prefix E that denotes an *entry_declaration*:
 27 Yields a value of the type Task_Id that identifies the task whose call is now being serviced. Use of this attribute is allowed only inside an *entry_body* or *accept_statement* corresponding to the *entry_declaration* denoted by E. See C.7.1.
- 28 S'Ceiling For every subtype S of a floating point type T:
 29 S'Ceiling denotes a function with the following specification:

$$\begin{array}{c} \text{function } S'\text{Ceiling } (X : T) \\ \text{return } T \end{array}$$

 30 The function yields the value $\lceil X \rceil$, i.e., the smallest (most negative) integral value greater than or equal to X. When X is zero, the result has the sign of X; a zero result otherwise has a negative sign when S'Signed_Zeros is True. See A.5.3.
- 31 S'Class For every subtype S of an untagged private type whose full view is tagged:
 32 Denotes the class-wide subtype corresponding to the full view of S. This attribute is allowed only from the beginning of the private part in which the full view is declared, until the declaration of the full view. After the full view, the Class attribute of the full view can be used. See 7.3.1.
- 33 S'Class For every subtype S of a tagged type T (specific or class-wide):
 34 S'Class denotes a subtype of the class-wide type (called T'Class in this International Standard) for the class rooted at T (or if S already denotes a class-wide subtype, then S'Class is the same as S).
 35 S'Class is unconstrained. However, if S is constrained, then the values of S'Class are only those that when converted to the type T belong to S. See 3.9.
- 36/1 X'Component_Size
 37 For a prefix X that denotes an array subtype or array object (after any implicit dereference):
 38 Denotes the size in bits of components of the type of X. The value of this attribute is of type universal_integer. See 13.3.
- 39 S'Compose For every subtype S of a floating point type T:
 40 S'Compose denotes a function with the following specification:

$$\begin{array}{c} \text{function } S'\text{Compose } (\textit{Fraction} : T; \\ \textit{Exponent} : \text{universal_integer}) \\ \text{return } T \end{array}$$

	Let v be the value $Fraction \cdot T\text{Machine_Radix}^{\text{Exponent}-k}$, where k is the normalized exponent of $Fraction$. If v is a machine number of the type T , or if $ v \geq T\text{Model_Small}$, the function yields v ; otherwise, it yields either one of the machine numbers of the type T adjacent to v . Constraint_Error is optionally raised if v is outside the base range of S . A zero result has the sign of $Fraction$ when S'Signed_Zeros is True. See A.5.3.	41
A'Constrained	For a prefix A that is of a discriminated type (after any implicit dereference):	42
	Yields the value True if A denotes a constant, a value, or a constrained variable, and False otherwise. See 3.7.2.	43
S'Copy_Sign	For every subtype S of a floating point type T :	44
	S'Copy_Sign denotes a function with the following specification:	45
	<code>function S'Copy_Sign (Value, Sign : T) return T</code>	46
	If the value of <i>Value</i> is nonzero, the function yields a result whose magnitude is that of <i>Value</i> and whose sign is that of <i>Sign</i> ; otherwise, it yields the value zero. Constraint_Error is optionally raised if the result is outside the base range of S. A zero result has the sign of <i>Sign</i> when S'Signed_Zeros is True. See A.5.3.	47
E'Count	For a prefix E that denotes an entry of a task or protected unit:	48
	Yields the number of calls presently queued on the entry E of the current instance of the unit. The value of this attribute is of the type universal_integer. See 9.9.	49
S'Definite	For a prefix S that denotes a formal indefinite subtype:	50/1
	S'Definite yields True if the actual subtype corresponding to S is definite; otherwise it yields False. The value of this attribute is of the predefined type Boolean. See 12.5.1.	51
S'Delta	For every fixed point subtype S:	52
	S'Delta denotes the <i>delta</i> of the fixed point subtype S. The value of this attribute is of the type universal_real. See 3.5.10.	53
S'Denorm	For every subtype S of a floating point type T :	54
	Yields the value True if every value expressible in the form $\pm \text{mantissa} \cdot T\text{Machine_Radix}^{T\text{Machine_Emin}}$	55
	where <i>mantissa</i> is a nonzero $T\text{Machine_Mantissa}$ -digit fraction in the number base $T\text{Machine_Radix}$, the first digit of which is zero, is a machine number (see 3.5.7) of the type T ; yields the value False otherwise. The value of this attribute is of the predefined type Boolean. See A.5.3.	
S'Digits	For every decimal fixed point subtype S:	56
	S'Digits denotes the <i>digits</i> of the decimal fixed point subtype S, which corresponds to the number of decimal digits that are representable in objects of the subtype. The value of this attribute is of the type universal_integer. See 3.5.10.	57
S'Digits	For every floating point subtype S:	58
	S'Digits denotes the requested decimal precision for the subtype S. The value of this attribute is of the type universal_integer. See 3.5.8.	59
S'Exponent	For every subtype S of a floating point type T :	60
	S'Exponent denotes a function with the following specification:	61
	<code>function S'Exponent (X : T) return universal_integer</code>	62

- 63 The function yields the normalized exponent of X . See A.5.3.
- 64 S'External_Tag
For every subtype S of a tagged type T (specific or class-wide):
- 65 S'External_Tag denotes an external string representation for S'Tag; it is of the predefined type String. External_Tag may be specified for a specific tagged type via an attribute_definition_clause; the expression of such a clause shall be static. The default external tag representation is implementation defined. See 3.9.2 and 13.13.2. See 13.3.
- 66/1 A'First For a prefix A that is of an array type (after any implicit dereference), or denotes a constrained array subtype:
- 67 A'First denotes the lower bound of the first index range; its type is the corresponding index type. See 3.6.2.
- 68 S'First For every scalar subtype S:
- 69 S'First denotes the lower bound of the range of S. The value of this attribute is of the type of S. See 3.5.
- 70/1 A'First(N) For a prefix A that is of an array type (after any implicit dereference), or denotes a constrained array subtype:
- 71 A'First(N) denotes the lower bound of the N-th index range; its type is the corresponding index type. See 3.6.2.
- 72 R.C'First_Bit
For a component C of a composite, non-array object R:
- 73/2 If the nondefault bit ordering applies to the composite type, and if a component_clause specifies the placement of C, denotes the value given for the first_bit of the component_clause; otherwise, denotes the offset, from the start of the first of the storage elements occupied by C, of the first bit occupied by C. This offset is measured in bits. The first bit of a storage element is numbered zero. The value of this attribute is of the type universal_integer. See 13.5.2.
- 74 S'Floor For every subtype S of a floating point type T:
- 75 S'Floor denotes a function with the following specification:
- ```
function S'Floor (X : T)
 return T
```
- 77                   The function yields the value  $\lfloor X \rfloor$ , i.e., the largest (most positive) integral value less than or equal to  $X$ . When  $X$  is zero, the result has the sign of  $X$ ; a zero result otherwise has a positive sign. See A.5.3.
- 78    S'Fore         For every fixed point subtype S:
- 79                   S'Fore yields the minimum number of characters needed before the decimal point for the decimal representation of any value of the subtype S, assuming that the representation does not include an exponent, but includes a one-character prefix that is either a minus sign or a space. (This minimum number does not include superfluous zeros or underlines, and is at least 2.) The value of this attribute is of the type universal\_integer. See 3.5.10.
- 80    S'Fraction     For every subtype S of a floating point type T:
- 81                   S'Fraction denotes a function with the following specification:
- ```
function S'Fraction (X : T)
    return T
```
- 83 The function yields the value $X \cdot T\text{Machine_Radix}^k$, where k is the normalized exponent of X . A zero result, which can only occur when X is zero, has the sign of X . See A.5.3.

T'Identity	For a prefix T that is of a task type (after any implicit dereference):	84
	Yields a value of the type Task_Id that identifies the task denoted by T. See C.7.1.	85
E'Identity	For a prefix E that denotes an exception:	86/1
	E'Identity returns the unique identity of the exception. The type of this attribute is Exception_Id. See 11.4.1.	87
S'Image	For every scalar subtype S:	88
	S'Image denotes a function with the following specification:	89
	<pre>function S'Image(Arg : S'Base) return String</pre>	90
	The function returns an image of the value of Arg as a String. See 3.5.	91/2
S'Class'Input	For every subtype S'Class of a class-wide type T'Class:	92
	S'Class'Input denotes a function with the following specification:	93
	<pre>function S'Class'Input(Stream : not null access Ada.Streams.Root_Stream_Type'Class) return T'Class</pre>	94/2
	First reads the external tag from Stream and determines the corresponding internal tag (by calling Tags.Descendant_Tag(String'Input(Stream), S'Tag) which might raise Tag_Error — see 3.9) and then dispatches to the subprogram denoted by the Input attribute of the specific type identified by the internal tag; returns that result. If the specific type identified by the internal tag is not covered by T'Class or is abstract, Constraint_Error is raised. See 13.13.2.	95/2
S'Input	For every subtype S of a specific type T:	96
	S'Input denotes a function with the following specification:	97
	<pre>function S'Input(Stream : not null access Ada.Streams.Root_Stream_Type'Class) return T</pre>	98/2
	S'Input reads and returns one value from Stream, using any bounds or discriminants written by a corresponding S'Output to determine how much to read. See 13.13.2.	99
A'Last	For a prefix A that is of an array type (after any implicit dereference), or denotes a constrained array subtype:	100/1
	A'Last denotes the upper bound of the first index range; its type is the corresponding index type. See 3.6.2.	101
S'Last	For every scalar subtype S:	102
	S'Last denotes the upper bound of the range of S. The value of this attribute is of the type of S. See 3.5.	103
A'Last(N)	For a prefix A that is of an array type (after any implicit dereference), or denotes a constrained array subtype:	104/1
	A'Last(N) denotes the upper bound of the N-th index range; its type is the corresponding index type. See 3.6.2.	105
R.C'Last_Bit	For a component C of a composite, non-array object R:	106
	If the nondefault bit ordering applies to the composite type, and if a component_clause specifies the placement of C, denotes the value given for the last_bit of the component_clause; otherwise, denotes the offset, from the start of the first of the storage	107/2

elements occupied by C, of the last bit occupied by C. This offset is measured in bits. The value of this attribute is of the type *universal_integer*. See 13.5.2.

108 S'Leading_Part

For every subtype S of a floating point type T:

S'Leading_Part denotes a function with the following specification:

```
110   function S'Leading_Part (X : T;
                           Radix_Digits : universal_integer)
                                return T
```

Let v be the value $T\text{Machine_Radix}^{k-\text{Radix_Digits}}$, where k is the normalized exponent of X. The function yields the value

- $\lfloor X/v \rfloor \cdot v$, when X is nonnegative and Radix_Digits is positive;
- $\lceil X/v \rceil \cdot v$, when X is negative and Radix_Digits is positive.

Constraint_Error is raised when Radix_Digits is zero or negative. A zero result, which can only occur when X is zero, has the sign of X. See A.5.3.

115/1 A'Length

For a prefix A that is of an array type (after any implicit dereference), or denotes a constrained array subtype:

A'Length denotes the number of values of the first index range (zero for a null range); its type is *universal_integer*. See 3.6.2.

117/1 A'Length(N)

For a prefix A that is of an array type (after any implicit dereference), or denotes a constrained array subtype:

A'Length(N) denotes the number of values of the N-th index range (zero for a null range); its type is *universal_integer*. See 3.6.2.

119 S'Machine

For every subtype S of a floating point type T:

S'Machine denotes a function with the following specification:

```
121   function S'Machine (X : T)
                            return T
```

If X is a machine number of the type T, the function yields X; otherwise, it yields the value obtained by rounding or truncating X to either one of the adjacent machine numbers of the type T. Constraint_Error is raised if rounding or truncating X to the precision of the machine numbers results in a value outside the base range of S. A zero result has the sign of X when S'Signed_Zeros is True. See A.5.3.

123 S'Machine_Emax

For every subtype S of a floating point type T:

Yields the largest (most positive) value of exponent such that every value expressible in the canonical form (for the type T), having a mantissa of TMachine_Mantissa digits, is a machine number (see 3.5.7) of the type T. This attribute yields a value of the type *universal_integer*. See A.5.3.

125 S'Machine_Emin

For every subtype S of a floating point type T:

Yields the smallest (most negative) value of exponent such that every value expressible in the canonical form (for the type T), having a mantissa of TMachine_Mantissa digits, is a machine number (see 3.5.7) of the type T. This attribute yields a value of the type *universal_integer*. See A.5.3.

127 S'Machine_Mantissa

For every subtype S of a floating point type T:

	Yields the largest value of p such that every value expressible in the canonical form (for the type T), having a p -digit <i>mantissa</i> and an <i>exponent</i> between $T\text{Machine_Emin}$ and $T\text{Machine_Emax}$, is a machine number (see 3.5.7) of the type T . This attribute yields a value of the type <i>universal_integer</i> . See A.5.3.	128
S'Machine_Overflows	For every subtype S of a fixed point type T :	129
	Yields the value True if overflow and divide-by-zero are detected and reported by raising Constraint_Error for every predefined operation that yields a result of the type T ; yields the value False otherwise. The value of this attribute is of the predefined type Boolean. See A.5.4.	130
S'Machine_Overflows	For every subtype S of a floating point type T :	131
	Yields the value True if overflow and divide-by-zero are detected and reported by raising Constraint_Error for every predefined operation that yields a result of the type T ; yields the value False otherwise. The value of this attribute is of the predefined type Boolean. See A.5.3.	132
S'Machine_Radix	For every subtype S of a fixed point type T :	133
	Yields the radix of the hardware representation of the type T . The value of this attribute is of the type <i>universal_integer</i> . See A.5.4.	134
S'Machine_Radix	For every subtype S of a floating point type T :	135
	Yields the radix of the hardware representation of the type T . The value of this attribute is of the type <i>universal_integer</i> . See A.5.3.	136
S'Machine_Rounding	For every subtype S of a floating point type T :	136.1/2
	S'Machine_Rounding denotes a function with the following specification:	136.2/2
	<code>function S'Machine_Rounding (X : T) return T</code>	136.3/2
	The function yields the integral value nearest to X . If X lies exactly halfway between two integers, one of those integers is returned, but which of them is returned is unspecified. A zero result has the sign of X when S'Signed_Zeros is True. This function provides access to the rounding behavior which is most efficient on the target processor. See A.5.3.	136.4/2
S'Machine_Rounds	For every subtype S of a fixed point type T :	137
	Yields the value True if rounding is performed on inexact results of every predefined operation that yields a result of the type T ; yields the value False otherwise. The value of this attribute is of the predefined type Boolean. See A.5.4.	138
S'Machine_Rounds	For every subtype S of a floating point type T :	139
	Yields the value True if rounding is performed on inexact results of every predefined operation that yields a result of the type T ; yields the value False otherwise. The value of this attribute is of the predefined type Boolean. See A.5.3.	140
S'Max	For every scalar subtype S:	141
	S'Max denotes a function with the following specification:	142

143 **function** S'Max(*Left*, *Right* : S'Base)
 return S'Base

144 The function returns the greater of the values of the two parameters. See 3.5.

145 S'Max_Size_In_Storage_Elements
 For every subtype S:

146/2 Denotes the maximum value for Size_In_Storage_Elements that could be requested by the implementation via Allocate for an access type whose designated subtype is S. For a type with access discriminants, if the implementation allocates space for a coextension in the same pool as that of the object having the access discriminant, then this accounts for any calls on Allocate that could be performed to provide space for such coextensions. The value of this attribute is of type *universal_integer*. See 13.11.1.

147 S'Min For every scalar subtype S:

148 S'Min denotes a function with the following specification:

149 **function** S'Min(*Left*, *Right* : S'Base)
 return S'Base

150 The function returns the lesser of the values of the two parameters. See 3.5.

150.1/2 S'Mod For every modular subtype S:

150.2/2 S'Mod denotes a function with the following specification:

150.3/2 **function** S'Mod (*Arg* : *universal_integer*)
 return S'Base

150.4/2 This function returns *Arg mod S'Modulus*, as a value of the type of S. See 3.5.4.

151 S'Model For every subtype S of a floating point type *T*:

152 S'Model denotes a function with the following specification:

153 **function** S'Model (*X* : *T*)
 return *T*

154 If the Numerics Annex is not supported, the meaning of this attribute is implementation defined; see G.2.2 for the definition that applies to implementations supporting the Numerics Annex. See A.5.3.

155 S'Model_Emin For every subtype S of a floating point type *T*:

156 If the Numerics Annex is not supported, this attribute yields an implementation defined value that is greater than or equal to the value of *TMachine_Emin*. See G.2.2 for further requirements that apply to implementations supporting the Numerics Annex. The value of this attribute is of the type *universal_integer*. See A.5.3.

157 S'Model_Epsilon For every subtype S of a floating point type *T*:

158 Yields the value $T\text{Machine_Radix}^{1 - T\text{Model_Mantissa}}$. The value of this attribute is of the type *universal_real*. See A.5.3.

159 S'Model_Mantissa For every subtype S of a floating point type *T*:

160 If the Numerics Annex is not supported, this attribute yields an implementation defined value that is greater than or equal to $\lceil d \cdot \log(10) / \log(T\text{Machine_Radix}) \rceil + 1$, where *d* is the requested decimal precision of *T*, and less than or equal to the value of *TMachine_Mantissa*. See G.2.2 for further requirements that apply to implementations

	supporting the Numerics Annex. The value of this attribute is of the type <i>universal_integer</i> . See A.5.3.	
S'Model_Small	For every subtype S of a floating point type T:	161
	Yields the value $T\text{Machine_Radix}^{T\text{Model_Emin} - 1}$. The value of this attribute is of the type <i>universal_real</i> . See A.5.3.	162
S'Modulus	For every modular subtype S:	163
	S'Modulus yields the modulus of the type of S, as a value of the type <i>universal_integer</i> . See 3.5.4.	164
S'Class'Output	For every subtype S'Class of a class-wide type T'Class:	165
	S'Class'Output denotes a procedure with the following specification:	166
	<pre>procedure S'Class'Output(Stream : not null access Ada.Streams.Root_Stream_Type'Class; Item : in T'Class)</pre>	167/2
	First writes the external tag of Item to Stream (by calling String'Output(Stream, Tags.External_Tag(Item'Tag)) — see 3.9) and then dispatches to the subprogram denoted by the Output attribute of the specific type identified by the tag. Tag_Error is raised if the tag of Item identifies a type declared at an accessibility level deeper than that of S. See 13.13.2.	168/2
S'Output	For every subtype S of a specific type T:	169
	S'Output denotes a procedure with the following specification:	170
	<pre>procedure S'Output(Stream : not null access Ada.Streams.Root_Stream_Type'Class; Item : in T)</pre>	171/2
	S'Output writes the value of Item to Stream, including any bounds or discriminants. See 13.13.2.	172
D'Partition_Id	For a prefix D that denotes a library-level declaration, excepting a declaration of or within a declared-pure library unit:	173/1
	Denotes a value of the type <i>universal_integer</i> that identifies the partition in which D was elaborated. If D denotes the declaration of a remote call interface library unit (see E.2.3) the given partition is the one where the body of D was elaborated. See E.1.	174
S'Pos	For every discrete subtype S:	175
	S'Pos denotes a function with the following specification:	176
	<pre>function S'Pos(Arg : S'Base) return universal_integer</pre>	177
	This function returns the position number of the value of Arg, as a value of type <i>universal_integer</i> . See 3.5.5.	178
R.C'Position	For a component C of a composite, non-array object R:	179
	If the nondefault bit ordering applies to the composite type, and if a component_clause specifies the placement of C, denotes the value given for the position of the component_clause; otherwise, denotes the same value as R.C'Address – R'Address. The value of this attribute is of the type <i>universal_integer</i> . See 13.5.2.	180/2
S'Pred	For every scalar subtype S:	181
	S'Pred denotes a function with the following specification:	182

183		<pre>function S'Pred(Arg : S'Base) return S'Base</pre>
184		For an enumeration type, the function returns the value whose position number is one less than that of the value of <i>Arg</i> ; <i>Constraint_Error</i> is raised if there is no such value of the type. For an integer type, the function returns the result of subtracting one from the value of <i>Arg</i> . For a fixed point type, the function returns the result of subtracting <i>small</i> from the value of <i>Arg</i> . For a floating point type, the function returns the machine number (as defined in 3.5.7) immediately below the value of <i>Arg</i> ; <i>Constraint_Error</i> is raised if there is no such machine number. See 3.5.
184.1/2	P'Priority	For a prefix P that denotes a protected object:
184.2/2		Denotes a non-aliased component of the protected object P. This component is of type System.Any_Priority and its value is the priority of P. P'Priority denotes a variable if and only if P denotes a variable. A reference to this attribute shall appear only within the body of P. See D.5.2.
185/1	A'Range	For a prefix A that is of an array type (after any implicit dereference), or denotes a constrained array subtype:
186		A'Range is equivalent to the range A'First .. A'Last, except that the prefix A is only evaluated once. See 3.6.2.
187	S'Range	For every scalar subtype S:
188		S'Range is equivalent to the range S'First .. S'Last. See 3.5.
189/1	A'Range(N)	For a prefix A that is of an array type (after any implicit dereference), or denotes a constrained array subtype:
190		A'Range(N) is equivalent to the range A'First(N) .. A'Last(N), except that the prefix A is only evaluated once. See 3.6.2.
191	S'Class'Read	For every subtype S'Class of a class-wide type T'Class:
192		S'Class'Read denotes a procedure with the following specification:
193/2		<pre>procedure S'Class'Read(Stream : not null access Ada.Streams.Root_Stream_Type'Class; Item : out T'Class)</pre>
194		Dispatches to the subprogram denoted by the Read attribute of the specific type identified by the tag of Item. See 13.13.2.
195	S'Read	For every subtype S of a specific type T:
196		S'Read denotes a procedure with the following specification:
197/2		<pre>procedure S'Read(Stream : not null access Ada.Streams.Root_Stream_Type'Class; Item : out T)</pre>
198		S'Read reads the value of Item from Stream. See 13.13.2.
199	S'Remainder	For every subtype S of a floating point type T:
200		S'Remainder denotes a function with the following specification:
201		<pre>function S'Remainder (X, Y : T) return T</pre>
202		For nonzero Y, let v be the value $X - n \cdot Y$, where n is the integer nearest to the exact value of X/Y ; if $ n - X/Y = 1/2$, then n is chosen to be even. If v is a machine number of the type T, the function yields v; otherwise, it yields zero. <i>Constraint_Error</i> is raised if Y is zero. A zero result has the sign of X when S'Signed_Zeros is True. See A.5.3.

S'Round	For every decimal fixed point subtype S:	203
	S'Round denotes a function with the following specification:	204
	<code>function S'Round (X : universal_real) return S'Base</code>	205
	The function returns the value obtained by rounding X (away from 0, if X is midway between two values of the type of S). See 3.5.10.	206
S'Rounding	For every subtype S of a floating point type T:	207
	S'Rounding denotes a function with the following specification:	208
	<code>function S'Rounding (X : T) return T</code>	209
	The function yields the integral value nearest to X, rounding away from zero if X lies exactly halfway between two integers. A zero result has the sign of X when S'Signed_Zeros is True. See A.5.3.	210
S'Safe_First	For every subtype S of a floating point type T:	211
	Yields the lower bound of the safe range (see 3.5.7) of the type T. If the Numerics Annex is not supported, the value of this attribute is implementation defined; see G.2.2 for the definition that applies to implementations supporting the Numerics Annex. The value of this attribute is of the type <i>universal_real</i> . See A.5.3.	212
S'Safe_Last	For every subtype S of a floating point type T:	213
	Yields the upper bound of the safe range (see 3.5.7) of the type T. If the Numerics Annex is not supported, the value of this attribute is implementation defined; see G.2.2 for the definition that applies to implementations supporting the Numerics Annex. The value of this attribute is of the type <i>universal_real</i> . See A.5.3.	214
S'Scale	For every decimal fixed point subtype S:	215
	S'Scale denotes the <i>scale</i> of the subtype S, defined as the value N such that S'Delta = 10.0**(-N). The scale indicates the position of the point relative to the rightmost significant digits of values of subtype S. The value of this attribute is of the type <i>universal_integer</i> . See 3.5.10.	216
S'Scaling	For every subtype S of a floating point type T:	217
	S'Scaling denotes a function with the following specification:	218
	<code>function S'Scaling (X : T; Adjustment : universal_integer) return T</code>	219
	Let v be the value $X \cdot T\text{Machine_Radix}^{\text{Adjustment}}$. If v is a machine number of the type T, or if $ v \geq T\text{Model_Small}$, the function yields v; otherwise, it yields either one of the machine numbers of the type T adjacent to v. Constraint_Error is optionally raised if v is outside the base range of S. A zero result has the sign of X when S'Signed_Zeros is True. See A.5.3.	220
S'Signed_Zeros	For every subtype S of a floating point type T:	221
	Yields the value True if the hardware representation for the type T has the capability of representing both positively and negatively signed zeros, these being generated and used by the predefined operations of the type T as specified in IEC 559:1989; yields the value False otherwise. The value of this attribute is of the predefined type Boolean. See A.5.3.	222
S'Size	For every subtype S:	223

- 224 If S is definite, denotes the size (in bits) that the implementation would choose for the
following objects of subtype S:
- 225 • A record component of subtype S when the record type is packed.
 - 226 • The formal parameter of an instance of `Unchecked_Conversion` that converts
from subtype S to some other subtype.
- 227 If S is indefinite, the meaning is implementation defined. The value of this attribute is of
the type `universal_integer`. See 13.3.
- 228/1 S'Size For a prefix X that denotes an object:
- 229 Denotes the size in bits of the representation of the object. The value of this attribute is of
the type `universal_integer`. See 13.3.
- 230 S'Small For every fixed point subtype S:
- 231 S'Small denotes the *small* of the type of S. The value of this attribute is of the type
`universal_real`. See 3.5.10.
- 232 S'Storage_Pool For every access-to-object subtype S:
- 233 Denotes the storage pool of the type of S. The type of this attribute is `Root_Storage_-
Pool'Class`. See 13.11.
- 234 S'Storage_Size For every access-to-object subtype S:
- 235 Yields the result of calling `Storage_Size(S'Storage_Pool)`, which is intended to be a
measure of the number of storage elements reserved for the pool. The type of this attribute
is `universal_integer`. See 13.11.
- 236/1 T'Storage_Size For a prefix T that denotes a task object (after any implicit dereference):
- 237 Denotes the number of storage elements reserved for the task. The value of this attribute is
of the type `universal_integer`. The Storage_Size includes the size of the task's stack, if any.
The language does not specify whether or not it includes other storage associated with the
task (such as the “task control block” used by some implementations.) See 13.3.
- 237.1/2 S'Stream_Size For every subtype S of an elementary type T:
- 237.2/2 Denotes the number of bits occupied in a stream by items of subtype S. Hence, the number
of stream elements required per item of elementary type T is:
- 237.3/2 `T'Stream_Size / Ada.Streams.Stream_Element'Size`
- 237.4/2 The value of this attribute is of type `universal_integer` and is a multiple of
`Stream_Element'Size`. See 13.13.2.
- 238 S'Succ For every scalar subtype S:
- 239 S'Succ denotes a function with the following specification:
- 240 `function S'Succ(Arg : S'Base)
 return S'Base`
- 241 For an enumeration type, the function returns the value whose position number is one more
than that of the value of Arg; `Constraint_Error` is raised if there is no such value of the type.
For an integer type, the function returns the result of adding one to the value of Arg. For a
fixed point type, the function returns the result of adding *small* to the value of Arg. For a
floating point type, the function returns the machine number (as defined in 3.5.7)

	immediately above the value of <i>Arg</i> ; Constraint_Error is raised if there is no such machine number. See 3.5.	
X'Tag	For a prefix X that is of a class-wide tagged type (after any implicit dereference):	242
	X'Tag denotes the tag of X. The value of this attribute is of type Tag. See 3.9.	243
S'Tag	For every subtype S of a tagged type T (specific or class-wide):	244
	S'Tag denotes the tag of the type T (or if T is class-wide, the tag of the root type of the corresponding class). The value of this attribute is of type Tag. See 3.9.	245
T'Terminated	For a prefix T that is of a task type (after any implicit dereference):	246
	Yields the value True if the task denoted by T is terminated, and False otherwise. The value of this attribute is of the predefined type Boolean. See 9.9.	247
S'Truncation	For every subtype S of a floating point type T:	248
	S'Truncation denotes a function with the following specification: <code>function S'Truncation (X : T) return T</code>	249
	The function yields the value $\lceil X \rceil$ when X is negative, and $\lfloor X \rfloor$ otherwise. A zero result has the sign of X when S'Signed_Zeros is True. See A.5.3.	251
S'Unbiased_Rounding	For every subtype S of a floating point type T:	252
	S'Unbiased_Rounding denotes a function with the following specification: <code>function S'Unbiased_Rounding (X : T) return T</code>	253
	The function yields the integral value nearest to X, rounding toward the even integer if X lies exactly halfway between two integers. A zero result has the sign of X when S'Signed_Zeros is True. See A.5.3.	255
X'Unchecked_Access	For a prefix X that denotes an aliased view of an object:	256
	All rules and semantics that apply to X'Access (see 3.10.2) apply also to X'Unchecked_Access, except that, for the purposes of accessibility rules and checks, it is as if X were declared immediately within a library package. See 13.10.	257
S'Val	For every discrete subtype S:	258
	S'Val denotes a function with the following specification: <code>function S'Val(Arg : universal_integer) return S'Base</code>	259
	This function returns a value of the type of S whose position number equals the value of Arg. See 3.5.5.	261
X'Valid	For a prefix X that denotes a scalar object (after any implicit dereference):	262
	Yields True if and only if the object denoted by X is normal and has a valid representation. The value of this attribute is of the predefined type Boolean. See 13.9.2.	263
S'Value	For every scalar subtype S:	264
	S'Value denotes a function with the following specification: <code>function S'Value(Arg : String) return S'Base</code>	265
		266

- 267 This function returns a value given an image of the value as a String, ignoring any leading or trailing spaces. See 3.5.
- 268/1 P'Version For a prefix P that statically denotes a program unit:
- 269 Yields a value of the predefined type String that identifies the version of the compilation unit that contains the declaration of the program unit. See E.3.
- 270 S'Wide_Image For every scalar subtype S:
- 271 S'Wide_Image denotes a function with the following specification:
- 272

```
function S'Wide_Image(Arg : S'Base)
                    return Wide_String
```
- 273/2 The function returns an image of the value of Arg as a Wide_String. See 3.5.
- 274 S'Wide_Value For every scalar subtype S:
- 275 S'Wide_Value denotes a function with the following specification:
- 276

```
function S'Wide_Value(Arg : Wide_String)
                    return S'Base
```
- 277 This function returns a value given an image of the value as a Wide_String, ignoring any leading or trailing spaces. See 3.5.
- 277.1/2 S'Wide_Wide_Image For every scalar subtype S:
- 277.2/2 S'Wide_Wide_Image denotes a function with the following specification:
- 277.3/2

```
function S'Wide_Wide_Image(Arg : S'Base)
                    return Wide_Wide_String
```
- 277.4/2 The function returns an *image* of the value of Arg, that is, a sequence of characters representing the value in display form. See 3.5.
- 277.5/2 S'Wide_Wide_Value For every scalar subtype S:
- 277.6/2 S'Wide_Wide_Value denotes a function with the following specification:
- 277.7/2

```
function S'Wide_Wide_Value(Arg : Wide_Wide_String)
                    return S'Base
```
- 277.8/2 This function returns a value given an image of the value as a Wide_Wide_String, ignoring any leading or trailing spaces. See 3.5.
- 277.9/2 S'Wide_Wide_Width For every scalar subtype S:
- 277.10/2 S'Wide_Wide_Width denotes the maximum length of a Wide_Wide_String returned by S'Wide_Wide_Image over all values of the subtype S. It denotes zero for a subtype that has a null range. Its type is universal_integer. See 3.5.
- 278 S'Wide_Width For every scalar subtype S:
- 279 S'Wide_Width denotes the maximum length of a Wide_String returned by S'Wide_Image over all values of the subtype S. It denotes zero for a subtype that has a null range. Its type is universal_integer. See 3.5.
- 280 S'Width For every scalar subtype S:

S'Width	denotes the maximum length of a String returned by S'Image over all values of the subtype S. It denotes zero for a subtype that has a null range. Its type is <i>universal_integer</i> . See 3.5.	281
S'Class'Write	For every subtype S'Class of a class-wide type T'Class:	282
	S'Class'Write denotes a procedure with the following specification:	283
	<pre>procedure S'Class'Write(Stream : not null access Ada.Streams.Root_Stream_Type'Class; Item : in T'Class)</pre>	284/2
	Dispatches to the subprogram denoted by the Write attribute of the specific type identified by the tag of Item. See 13.13.2.	285
S'Write	For every subtype S of a specific type T:	286
	S'Write denotes a procedure with the following specification:	287
	<pre>procedure S'Write(Stream : not null access Ada.Streams.Root_Stream_Type'Class; Item : in T)</pre>	288/2
	S'Write writes the value of Item to Stream. See 13.13.2.	289

Annex L (informative) Language-Defined Pragmas

This Annex summarizes the definitions given elsewhere of the language-defined pragmas.	1
pragma All_Calls_Remote[(<i>library_unit_name</i>)]; — See E.2.3.	2
pragma Assert([Check =>] <i>boolean_expression</i>[, [Message =>] <i>string_expression</i>]); — See 11.4.2.	2.1/2
pragma Assertion_Policy(<i>policy_identifier</i>); — See 11.4.2.	2.2/2
pragma Asynchronous(<i>local_name</i>); — See E.4.1.	3
pragma Atomic(<i>local_name</i>); — See C.6.	4
pragma Atomic_Components(<i>array_local_name</i>); — See C.6.	5
pragma Attach_Handler(<i>handler_name</i>, <i>expression</i>); — See C.3.1.	6
pragma Controlled(<i>first_subtype_local_name</i>); — See 13.11.3.	7
pragma Convention([Convention =>] <i>convention_identifier</i>,[Entity =>] <i>local_name</i>); — See B.1.	8
pragma Detect_Blocking; — See H.5.	8.1/2
pragma Discard_Names([(On =>] <i>local_name</i>)]; — See C.5.	9
pragma Elaborate(<i>library_unit_name</i>{, <i>library_unit_name</i>}); — See 10.2.1.	10
pragma Elaborate_All(<i>library_unit_name</i>{, <i>library_unit_name</i>}); — See 10.2.1.	11
pragma Elaborate_Body(<i>library_unit_name</i>); — See 10.2.1.	12
pragma Export([Convention =>] <i>convention_identifier</i> , [Entity =>] <i>local_name</i> [, [External_Name =>] <i>string_expression</i>] [, [Link_Name =>] <i>string_expression</i>]); — See B.1.	13
pragma Import([Convention =>] <i>convention_identifier</i> , [Entity =>] <i>local_name</i> [, [External_Name =>] <i>string_expression</i>] [, [Link_Name =>] <i>string_expression</i>]); — See B.1.	14
pragma Inline(<i>name</i> {, <i>name</i>}); — See 6.3.2.	15
pragma Inspection_Point([(<i>object_name</i> {, <i>object_name</i>}))]; — See H.3.2.	16
pragma Interrupt_Handler(<i>handler_name</i>); — See C.3.1.	17
pragma Interrupt_Priority([<i>expression</i>]); — See D.1.	18
pragma Linker_Options(<i>string_expression</i>); — See B.1.	19
pragma List(<i>identifier</i>); — See 2.8.	20
pragma Locking_Policy(<i>policy_identifier</i>); — See D.3.	21
pragma No_Return(<i>procedure_local_name</i>{, <i>procedure_local_name</i>}); — See 6.5.1.	21.1/2
pragma Normalize_Scalars; — See H.1.	22

- 23 **pragma** Optimize(*identifier*); — See 2.8.
- 24 **pragma** Pack(*first_subtype_local_name*); — See 13.2.
- 25 **pragma** Page; — See 2.8.
- 25.1/2 **pragma** Partition_Elaboration_Policy (*policy_identifier*); — See H.6.
- 25.2/2 **pragma** Preelaborable_Initialization(*direct_name*); — See 10.2.1.
- 26 **pragma** Preelaborate[*(library_unit_name)*]; — See 10.2.1.
- 27 **pragma** Priority(*expression*); — See D.1.
- 27.1/2 **pragma** Priority_Specific_Dispatching (*policy_identifier*, *first_priority_expression*, *last_priority_expression*); — See D.2.2.
- 27.2/2 **pragma** Profile (*profile_identifier* {, *profile pragma_argument_association*}); — See D.13.
- 28 **pragma** Pure[*(library_unit_name)*]; — See 10.2.1.
- 29 **pragma** Queuing_Policy(*policy_identifier*); — See D.4.
- 29.1/2 **pragma** Relative_Deadline (*relative_deadline_expression*); — See D.2.6.
- 30 **pragma** Remote_Call_Interface[*(library_unit_name)*]; — See E.2.3.
- 31 **pragma** Remote_Types[*(library_unit_name)*]; — See E.2.2.
- 32 **pragma** Restrictions(*restriction* {, *restriction*}); — See 13.12.
- 33 **pragma** Reviewable; — See H.3.1.
- 34 **pragma** Shared_Passive[*(library_unit_name)*]; — See E.2.1.
- 35 **pragma** Storage_Size(*expression*); — See 13.3.
- 36 **pragma** Suppress(*identifier*); — See 11.5.
- 37 **pragma** Task_Dispatching_Policy(*policy_identifier*); — See D.2.2.
- 37.1/2 **pragma** Unchecked_Union (*first_subtype_local_name*); — See B.3.3.
- 37.2/2 **pragma** Unsuppress(*identifier*); — See 11.5.
- 38 **pragma** Volatile(*local_name*); — See C.6.
- 39 **pragma** Volatile_Components(*array_local_name*); — See C.6.

Annex M (informative) Summary of Documentation Requirements

The Ada language allows for certain target machine dependences in a controlled manner. Each Ada implementation must document many characteristics and properties of the target system. This International Standard contains specific documentation requirements. In addition, many characteristics that require documentation are identified throughout this International Standard as being implementation defined. Finally, this International Standard requires documentation of whether implementation advice is followed. The following clauses provide summaries of these documentation requirements.

M.1 Specific Documentation Requirements

In addition to implementation-defined characteristics, each Ada implementation must document various properties of the implementation:

- The behavior of implementations in implementation-defined situations shall be documented — see M.2, “Implementation-Defined Characteristics” for a listing. See 1.1.3(19). 1/2
- The set of values that a user-defined Allocate procedure needs to accept for the Alignment parameter. How the standard storage pool is chosen, and how storage is allocated by standard storage pools. See 13.11(22). 3/2
- The algorithm used for random number generation, including a description of its period. See A.5.2(44). 4/2
- The minimum time interval between calls to the time-dependent Reset procedure that is guaranteed to initiate different random number sequences. See A.5.2(45). 5/2
- The conditions under which Io_Exceptions.Name_Error, Io_Exceptions.Use_Error, and Io_Exceptions.Device_Error are propagated. See A.13(15). 6/2
- The behavior of package Environment_Variables when environment variables are changed by external mechanisms. See A.17(30/2). 7/2
- The overhead of calling machine-code or intrinsic subprograms. See C.1(6). 8/2
- The types and attributes used in machine code insertions. See C.1(7). 9/2
- The subprogram calling conventions for all supported convention identifiers. See C.1(8). 10/2
- The mapping between the Link_Name or Ada designator and the external link name. See C.1(9). 11/2
- The treatment of interrupts. See C.3(22). 12/2
- The metrics for interrupt handlers. See C.3.1(16). 13/2
- If the Ceiling_Locking policy is in effect, the default ceiling priority for a protected object that contains an interrupt handler pragma. See C.3.2(24/2). 14/2
- Any circumstances when the elaboration of a preelaborated package causes code to be executed. See C.4(12). 15/2
- Whether a partition can be restarted without reloading. See C.4(13). 16/2
- The effect of calling Current_Task from an entry body or interrupt handler. See C.7.1(19). 17/2

- 18/2 • For package Task_Attributes, limits on the number and size of task attributes, and how to configure any limits. See C.7.2(19).
- 19/2 • The metrics for the Task_Attributes package. See C.7.2(27).
- 20/2 • The details of the configuration used to generate the values of all metrics. See D(2).
- 21/2 • The maximum priority inversion a user task can experience from the implementation. See D.2.3(12/2).
- 22/2 • The amount of time that a task can be preempted for processing on behalf of lower-priority tasks. See D.2.3(13/2).
- 23/2 • The quantum values supported for round robin dispatching. See D.2.5(16/2).
- 24/2 • The accuracy of the detection of the exhaustion of the budget of a task for round robin dispatching. See D.2.5(17/2).
- 25/2 • Any conditions that cause the completion of the setting of the deadline of a task to be delayed for a multiprocessor. See D.2.6(32/2).
- 26/2 • Any conditions that cause the completion of the setting of the priority of a task to be delayed for a multiprocessor. See D.5.1(12.1/2).
- 27/2 • The metrics for Set_Priority. See D.5.1(14).
- 28/2 • The metrics for setting the priority of a protected object. See D.5.2(10).
- 29/2 • On a multiprocessor, any conditions that cause the completion of an aborted construct to be delayed later than what is specified for a single processor. See D.6(3).
- 30/2 • The metrics for aborts. See D.6(8).
- 31/2 • The values of Time_First, Time_Last, Time_Span_First, Time_Span_Last, Time_Span_Unit, and Tick for package Real_Time. See D.8(33).
- 32/2 • The properties of the underlying time base used in package Real_Time. See D.8(34).
- 33/2 • Any synchronization of package Real_Time with external time references. See D.8(35).
- 34/2 • Any aspects of the external environment that could interfere with package Real_Time. See D.8(36/1).
- 35/2 • The metrics for package Real_Time. See D.8(45).
- 36/2 • The minimum value of the delay expression of a delay_relative_statement that causes a task to actually be blocked. See D.9(7).
- 37/2 • The minimum difference between the value of the delay expression of a delay_until_statement and the value of Real_Time.Clock, that causes the task to actually be blocked. See D.9(8).
- 38/2 • The metrics for delay statements. See D.9(13).
- 39/2 • The upper bound on the duration of interrupt blocking caused by the implementation. See D.12(5).
- 40/2 • The metrics for entry-less protected objects. See D.12(12).
- 41/2 • The values of CPU_Time_First, CPU_Time_Last, CPU_Time_Unit, and CPU_Tick of package Execution_Time. See D.14(21/2).
- 42/2 • The properties of the mechanism used to implement package Execution_Time. See D.14(22/2).
- 43/2 • The metrics for execution time. See D.14(27).
- 44/2 • The metrics for timing events. See D.15(24).

- Whether the RPC-receiver is invoked from concurrent tasks, and if so, the number of such tasks. See E.5(25). 45/2
- Any techniques used to reduce cancellation errors in Numerics.Generic_Real_Arrays shall be documented. See G.3.1(86/2). 46/2
- Any techniques used to reduce cancellation errors in Numerics.Generic_Complex_Arrays shall be documented. See G.3.2(155/2). 47/2
- If a **pragma** Normalize_Scalars applies, the implicit initial values of scalar subtypes shall be documented. Such a value should be an invalid representation when possible; any cases when it is not shall be documented. See H.1(5/2). 48/2
- The range of effects for each bounded error and each unspecified effect. If the effects of a given erroneous construct are constrained, the constraints shall be documented. See H.2(1). 49/2
- For each inspection point, a mapping between each inspectable object and the machine resources where the object's value can be obtained shall be provided. See H.3.2(8). 50/2
- If a **pragma** Restrictions(No_Exceptions) is specified, the effects of all constructs where language-defined checks are still performed. See H.4(25). 51/2
- The interrupts to which a task entry may be attached. See J.7.1(12). 52/2
- The type of entry call invoked for an interrupt entry. See J.7.1(13). 53/2

M.2 Implementation-Defined Characteristics

The Ada language allows for certain machine dependences in a controlled manner. Each Ada implementation must document all implementation-defined characteristics: 1/2

- Whether or not each recommendation given in Implementation Advice is followed — see M.3, “Implementation Advice” for a listing. See 1.1.2(37). 2/2
- Capacity limitations of the implementation. See 1.1.3(3). 3
- Variations from the standard that are impractical to avoid given the implementation's execution environment. See 1.1.3(6). 4
- Which **code_statements** cause external interactions. See 1.1.3(10). 5
- The semantics of an Ada program whose text is not in Normalization Form KC. See 2.1(4.1/2). 5.1/2
- The coded representation for the text of an Ada program. See 2.1(4/2). 6
- *This paragraph was deleted.* 7/2
- The representation for an end of line. See 2.2(2/2). 8
- Maximum supported line length and lexical element length. See 2.2(14). 9
- Implementation-defined pragmas. See 2.8(14). 10
- Effect of **pragma** Optimize. See 2.8(27). 11
- The sequence of characters of the value returned by S'Wide_Image when some of the graphic characters of S'Wide_Wide_Image are not defined in Wide_Character. See 3.5(30/2). 11.1/2
- The sequence of characters of the value returned by S'Image when some of the graphic characters of S'Wide_Wide_Image are not defined in Character. See 3.5(37/2). 12/2
- The predefined integer types declared in Standard. See 3.5.4(25). 13
- Any nonstandard integer types and the operators defined for them. See 3.5.4(26). 14

- 15
 - Any nonstandard real types and the operators defined for them. See 3.5.6(8).
- 16
 - What combinations of requested decimal precision and range are supported for floating point types. See 3.5.7(7).
- 17
 - The predefined floating point types declared in Standard. See 3.5.7(16).
- 18
 - The *small* of an ordinary fixed point type. See 3.5.9(8/2).
- 19
 - What combinations of *small*, range, and *digits* are supported for fixed point types. See 3.5.9(10).
- 19.1/2
 - The sequence of characters of the value returned by Tags.Expanded_Name (respectively, Tags.Wide_Expanded_Name) when some of the graphic characters of Tags.Wide_Wide_Expanded_Name are not defined in Character (respectively, Wide_Character). See 3.9(10.1/2).
- 20/2
 - The result of Tags.Wide_Wide_Expanded_Name for types declared within an unnamed block_statement. See 3.9(10).
- 21
 - Implementation-defined attributes. See 4.1.4(12/1).
- 21.1/2
 - Rounding of real static expressions which are exactly half-way between two machine numbers. See 4.9(38/2).
- 22
 - Any implementation-defined time types. See 9.6(6).
- 23
 - The time base associated with relative delays. See 9.6(20).
- 24
 - The time base of the type Calendar.Time. See 9.6(23).
- 25/2
 - The time zone used for package Calendar operations. See 9.6(24/2).
- 26
 - Any limit on delay_until_statements of select_statements. See 9.6(29).
- 26.1/2
 - The result of Calendar.Formatting.Image if its argument represents more than 100 hours. See 9.6.1(86/2).
- 27
 - Whether or not two nonoverlapping parts of a composite object are independently addressable, in the case where packing, record layout, or Component_Size is specified for the object. See 9.10(1).
- 28
 - The representation for a compilation. See 10.1(2).
- 29
 - Any restrictions on compilations that contain multiple compilation_units. See 10.1(4).
- 29.1/2
 - The mechanisms for adding a compilation unit mentioned in a limited_with_clause to an environment. See 10.1.4(3/2).
- 30
 - The mechanisms for creating an environment and for adding and replacing compilation units. See 10.1.4(3/2).
- 31
 - The implementation-defined means, if any, of specifying which compilation units are needed by a given compilation unit. See 10.2(2).
- 32
 - The manner of explicitly assigning library units to a partition. See 10.2(2).
- 33
 - The manner of designating the main subprogram of a partition. See 10.2(7).
- 34
 - The order of elaboration of library_items. See 10.2(18).
- 35
 - Parameter passing and function return for the main subprogram. See 10.2(21).
- 36
 - The mechanisms for building and running partitions. See 10.2(24).
- 37
 - The details of program execution, including program termination. See 10.2(25).
- 38
 - The semantics of any nonactive partitions supported by the implementation. See 10.2(28).

• The information returned by <code>Exception_Message</code> . See 11.4.1(10.1/2).	39
• The sequence of characters of the value returned by <code>Exceptions.Exception_Name</code> (respectively, <code>Exceptions.Wide_Exception_Name</code>) when some of the graphic characters of <code>Exceptions.Wide_Wide_Exception_Name</code> are not defined in <code>Character</code> (respectively, <code>Wide_Character</code>). See 11.4.1(12.1/2).	39.1/2
• The result of <code>Exceptions.Wide_Wide_Exception_Name</code> for exceptions declared within an unnamed <code>block_statement</code> . See 11.4.1(12).	40/2
• The information returned by <code>Exception_Information</code> . See 11.4.1(13/2).	41
• Implementation-defined <i>policy_identifiers</i> allowed in a <code>pragma Assertion_Policy</code> . See 11.4.2(9/2).	41.1/2
• The default assertion policy. See 11.4.2(10/2).	41.2/2
• Existence and meaning of second parameter of <code>pragma Unsuppress</code> . See 11.5(27.1/2).	41.3/2
• Implementation-defined check names. See 11.5(27).	42
• The cases that cause conflicts between the representation of the ancestors of a <code>type_declaration</code> . See 13.1(13.1/2).	42.1/2
• Any restrictions placed upon representation items. See 13.1(20).	43
• The interpretation of each aspect of representation. See 13.1(20).	44
• The set of machine scalars. See 13.3(8.1/2).	44.1/2
• The meaning of <code>Size</code> for indefinite subtypes. See 13.3(48).	45
• The default external representation for a type tag. See 13.3(75/1).	46
• What determines whether a compilation unit is the same in two different partitions. See 13.3(76).	47
• Implementation-defined components. See 13.5.1(15).	48
• If <code>Word_Size = Storage_Unit</code> , the default bit ordering. See 13.5.3(5).	49
• The contents of the visible part of package <code>System</code> . See 13.7(2).	50/2
• The range of <code>Storage_Elements.Storage_Offset</code> , the modulus of <code>Storage_Elements.Storage_Element</code> , and the declaration of <code>Storage_Elements.Integer_Address</code> . See 13.7.1(11).	50.1/2
• The contents of the visible part of package <code>System.Machine_Code</code> , and the meaning of <code>code_statements</code> . See 13.8(7).	51
• The effect of unchecked conversion for instances with nonscalar result types whose effect is not defined by the language. See 13.9(11).	52/2
• The result of unchecked conversion for instances with scalar result types whose result is not defined by the language. See 13.9(11).	52.1/2
• Whether or not the implementation provides user-accessible names for the standard pool type(s). See 13.11(17).	53
• <i>This paragraph was deleted.</i>	54/2
• The meaning of <code>Storage_Size</code> when neither the <code>Storage_Size</code> nor the <code>Storage_Pool</code> is specified for an access type. See 13.11(18).	55/2
• <i>This paragraph was deleted.</i>	56/2
• The set of restrictions allowed in a <code>pragma Restrictions</code> . See 13.12(7/2).	57/2

- 58 • The consequences of violating limitations on Restrictions pragmas. See 13.12(9).
- 59/2 • The contents of the stream elements read and written by the Read and Write attributes of elementary types. See 13.13.2(9).
- 60 • The names and characteristics of the numeric subtypes declared in the visible part of package Standard. See A.1(3).
- 60.1/2 • The values returned by Strings.Hash. See A.4.9(3/2).
- 61 • The accuracy actually achieved by the elementary functions. See A.5.1(1).
- 62 • The sign of a zero result from some of the operators or functions in Numerics.Generic_Elementary_Functions, when Float_Type'Signed_Zeros is True. See A.5.1(46).
- 63 • The value of Numerics.Discrete_Random.Max_Image_Width. See A.5.2(27).
- 64 • The value of Numerics.Float_Random.Max_Image_Width. See A.5.2(27).
- 65/2 • *This paragraph was deleted.*
- 66 • The string representation of a random number generator's state. See A.5.2(38).
- 67/2 • *This paragraph was deleted.*
- 68 • The values of the Model_Mantissa, Model_Emin, Model_Epsilon, Model, Safe_First, and Safe_Last attributes, if the Numerics Annex is not supported. See A.5.3(72).
- 69/2 • *This paragraph was deleted.*
- 70 • The value of Buffer_Size in Storage_IO. See A.9(10).
- 71/2 • The external files associated with the standard input, standard output, and standard error files. See A.10(5).
- 72 • The accuracy of the value produced by Put. See A.10.9(36).
- 72.1/1 • Current size for a stream file for which positioning is not supported. See A.12.1(1.1/1).
- 73/2 • The meaning of Argument_Count, Argument, and Command_Name for package Command_Line. The bounds of type Command_Line.Exit_Status. See A.15(1).
- 73.1/2 • The interpretation of file names and directory names. See A.16(46/2).
- 73.2/2 • The maximum value for a file size in Directories. See A.16(87/2).
- 73.3/2 • The result for Directories.Size for a directory or special file. See A.16(93/2).
- 73.4/2 • The result for Directories.Modification_Time for a directory or special file. See A.16(95/2).
- 73.5/2 • The interpretation of a non-null search pattern in Directories. See A.16(104/2).
- 73.6/2 • The results of a Directories search if the contents of the directory are altered while a search is in progress. See A.16(110/2).
- 73.7/2 • The definition and meaning of an environment variable. See A.17(1/2).
- 73.8/2 • The circumstances where an environment variable cannot be defined. See A.17(16/2).
- 73.9/2 • Environment names for which Set has the effect of Clear. See A.17(17/2).
- 73.10/2 • The value of Containers.Hash_Type'Modulus. The value of Containers.Count_Type'Last. See A.18.1(7/2).
- 74 • Implementation-defined convention names. See B.1(11).

• The manner of choosing link names when neither the link name nor the address of an imported or exported entity is specified. See B.1(36).	75
• The meaning of link names. See B.1(36).	76
• The effect of pragma Linker_Options. See B.1(37).	77
• The contents of the visible part of package Interfaces and its language-defined descendants. See B.2(1).	78
• Implementation-defined children of package Interfaces. See B.2(11).	79/2
• The definitions of certain types and constants in Interfaces.C. See B.3(41).	79.1/2
• The types Floating, Long_Floating, Binary, Long_Binary, Decimal_Element, and COBOL_Character; and the initializations of the variables Ada_To_COBOL and COBOL_To_Ada, in Interfaces.COBOL. See B.4(50).	80/1
• The types Fortran_Integer, Real, Double_Precision, and Character_Set in Interfaces.Fortran. See B.5(17).	80.1/1
• Implementation-defined intrinsic subprograms. See C.1(1).	81/2
• <i>This paragraph was deleted.</i>	82/2
• <i>This paragraph was deleted.</i>	83/2
• Any restrictions on a protected procedure or its containing type when a pragma Attach_Handler or Interrupt_Handler applies. See C.3.1(17).	83.1/2
• Any other forms of interrupt handler supported by the Attach_Handler and Interrupt_Handler pragmas. See C.3.1(19).	83.2/2
• <i>This paragraph was deleted.</i>	84/2
• The semantics of pragma Discard_Names . See C.5(7).	85
• The result of the Task_Identification.Image attribute. See C.7.1(7).	86
• The value of Current_Task when in a protected entry, interrupt handler, or finalization of a task attribute. See C.7.1(17/2).	87/2
• <i>This paragraph was deleted.</i>	88/2
• Granularity of locking for Task_Attributes . See C.7.2(16/1).	88.1/1
• <i>This paragraph was deleted.</i>	89/2
• <i>This paragraph was deleted.</i>	90/2
• The declarations of Any_Priority and Priority . See D.1(11).	91
• Implementation-defined execution resources. See D.1(15).	92
• Whether, on a multiprocessor, a task that is waiting for access to a protected object keeps its processor busy. See D.2.1(3).	93
• The effect of implementation-defined execution resources on task dispatching. See D.2.1(9/2).	94/2
• <i>This paragraph was deleted.</i>	95/2
• <i>This paragraph was deleted.</i>	96/2
• Implementation defined task dispatching policies. See D.2.2(18).	97/2
• The value of Default_Quantum in Dispatching.Round_Robin . See D.2.5(4).	97.1/2
• Implementation-defined policy_identifiers allowed in a pragma Locking_Policy . See D.3(4).	98

- 98.1/2 • The locking policy if no Locking_Policy pragma applies to any unit of a partition. See D.3(6).
- 99 • Default ceiling priorities. See D.3(10/2).
- 100 • The ceiling of any protected object used internally by the implementation. See D.3(16).
- 101 • Implementation-defined queuing policies. See D.4(1/1).
- 102/2 • *This paragraph was deleted.*
- 103 • Any operations that implicitly require heap storage allocation. See D.7(8).
- 103.1/2 • When restriction No_Task_Termination applies to a partition, what happens when a task terminates. See D.7(15.1/2).
- 103.2/2 • The behavior when restriction Max_Storage_At_Blocking is violated. See D.7(17/1).
- 103.3/2 • The behavior when restriction Max_Asynchronous_Select_Nesting is violated. See D.7(18/1).
- 103.4/2 • The behavior when restriction Max_Tasks is violated. See D.7(19).
- 104/2 • Whether the use of pragma Restrictions results in a reduction in program code or data size or execution time. See D.7(20).
- 105/2 • *This paragraph was deleted.*
- 106/2 • *This paragraph was deleted.*
- 107/2 • *This paragraph was deleted.*
- 108 • The means for creating and executing distributed programs. See E(5).
- 109 • Any events that can result in a partition becoming inaccessible. See E.1(7).
- 110 • The scheduling policies, treatment of priorities, and management of shared resources between partitions in certain cases. See E.1(11).
- 111/1 • *This paragraph was deleted.*
- 112 • Whether the execution of the remote subprogram is immediately aborted as a result of cancellation. See E.4(13).
- 112.1/2 • The range of type System.RPC.Partition_Id. See E.5(14).
- 113/2 • *This paragraph was deleted.*
- 114 • Implementation-defined interfaces in the PCS. See E.5(26).
- 115 • The values of named numbers in the package Decimal. See F.2(7).
- 116 • The value of Max_Picture_Length in the package Text_IO.Editing See F.3.3(16).
- 117 • The value of Max_Picture_Length in the package Wide_Text_IO.Editing See F.3.4(5).
- 117.1/2 • The value of Max_Picture_Length in the package Wide_Wide_Text_IO.Editing See F.3.5(5).
- 118 • The accuracy actually achieved by the complex elementary functions and by other complex arithmetic operations. See G.1(1).
- 119 • The sign of a zero result (or a component thereof) from any operator or function in Numerics.Generic_Complex_Types, when Real'Signed_Zeros is True. See G.1.1(53).
- 120 • The sign of a zero result (or a component thereof) from any operator or function in Numerics.Generic_Complex_Elementary_Functions, when Complex_Types.Real'Signed_Zeros is True. See G.1.2(45).
- 121 • Whether the strict mode or the relaxed mode is the default. See G.2(2).

- The result interval in certain cases of fixed-to-float conversion. See G.2.1(10). 122
- The result of a floating point arithmetic operation in overflow situations, when the Machine_Overflows attribute of the result type is False. See G.2.1(13). 123
- The result interval for division (or exponentiation by a negative exponent), when the floating point hardware implements division as multiplication by a reciprocal. See G.2.1(16). 124
- The definition of *close result set*, which determines the accuracy of certain fixed point multiplications and divisions. See G.2.3(5). 125
- Conditions on a *universal_real* operand of a fixed point multiplication or division for which the result shall be in the *perfect result set*. See G.2.3(22). 126
- The result of a fixed point arithmetic operation in overflow situations, when the Machine_Overflows attribute of the result type is False. See G.2.3(27). 127
- The result of an elementary function reference in overflow situations, when the Machine_Overflows attribute of the result type is False. See G.2.4(4). 128
- The accuracy of certain elementary functions for parameters beyond the angle threshold. See G.2.4(10). 129
- The value of the *angle threshold*, within which certain elementary functions, complex arithmetic operations, and complex elementary functions yield results conforming to a maximum relative error bound. See G.2.4(10). 130
- The result of a complex arithmetic operation or complex elementary function reference in overflow situations, when the Machine_Overflows attribute of the corresponding real type is False. See G.2.6(5). 131
- The accuracy of certain complex arithmetic operations and certain complex elementary functions for parameters (or components thereof) beyond the angle threshold. See G.2.6(8). 132
- The accuracy requirements for the subprograms Solve, Inverse, Determinant, Eigenvalues and Eigensystem for type Real_Matrix. See G.3.1(81/2). 132.1/2
- The accuracy requirements for the subprograms Solve, Inverse, Determinant, Eigenvalues and Eigensystem for type Complex_Matrix. See G.3.2(149/2). 132.2/2
- *This paragraph was deleted.* 133/2
- *This paragraph was deleted.* 134/2
- *This paragraph was deleted.* 135/2
- *This paragraph was deleted.* 136/2
- Implementation-defined *policy_identifiers* allowed in a *pragma Partition_Elaboration_Policy*. See H.6(4/2). 136.1/2

M.3 Implementation Advice

This International Standard sometimes gives advice about handling certain target machine dependences. 1/2
Each Ada implementation must document whether that advice is followed:

- Program_Error should be raised when an unsupported Specialized Needs Annex feature is used at run time. See 1.1.3(20). 2/2
- Implementation-defined extensions to the functionality of a language-defined library unit should be provided by adding children to the library unit. See 1.1.3(21). 3/2

- 4/2 • If a bounded error or erroneous execution is detected, Program_Error should be raised. See 1.1.5(12).
- 5/2 • Implementation-defined pragmas should have no semantic effect for error-free programs. See 2.8(16).
- 6/2 • Implementation-defined pragmas should not make an illegal program legal, unless they complete a declaration or configure the library_items in an environment. See 2.8(19).
- 7/2 • Long_Integer should be declared in Standard if the target supports 32-bit arithmetic. No other named integer subtypes should be declared in Standard. See 3.5.4(28).
- 8/2 • For a two's complement target, modular types with a binary modulus up to System.Max_Int*2+2 should be supported. A nonbinary modulus up to Integer'Last should be supported. See 3.5.4(29).
- 9/2 • Program_Error should be raised for the evaluation of S'Pos for an enumeration type, if the value of the operand does not correspond to the internal code for any enumeration literal of the type. See 3.5.5(8).
- 10/2 • Long_Float should be declared in Standard if the target supports 11 or more digits of precision. No other named float subtypes should be declared in Standard. See 3.5.7(17).
- 11/2 • Multidimensional arrays should be represented in row-major order, unless the array has convention Fortran. See 3.6.2(11).
- 12/2 • Tags.Internal_Tag should return the tag of a type whose innermost master is the master of the point of the function call.. See 3.9(26.1/2).
- 13/2 • For a real static expression with a non-formal type that is not part of a larger static expression should be rounded the same as the target system. See 4.9(38.1/2).
- 14/2 • The value of Duration'Small should be no greater than 100 microseconds. See 9.6(30).
- 15/2 • The time base for delay_relative_statements should be monotonic. See 9.6(31).
- 16/2 • Leap seconds should be supported if the target system supports them. Otherwise, operations in Calendar.Formatting should return results consistent with no leap seconds. See 9.6.1(89/2).
- 17/2 • When applied to a generic unit, a program unit pragma that is not a library unit pragma should apply to each instance of the generic unit for which there is not an overriding pragma applied directly to the instance. See 10.1.5(10/1).
- 18/2 • A type declared in a preelaborated package should have the same representation in every elaboration of a given version of the package. See 10.2.1(12).
- 19/2 • Exception_Message by default should be short, provide information useful for debugging, and should not include the Exception_Name. See 11.4.1(19).
- 20/2 • Exception_Information should provide information useful for debugging, and should include the Exception_Name and Exception_Message. See 11.4.1(19).
- 21/2 • Code executed for checks that have been suppressed should be minimized. See 11.5(28).
- 22/2 • The recommended level of support for all representation items should be followed. See 13.1(28/2).
- 23/2 • Storage allocated to objects of a packed type should be minimized. See 13.2(6).
- 24/2 • The recommended level of support for pragma Pack should be followed. See 13.2(9).
- 25/2 • For an array X, X'Address should point at the first component of the array rather than the array bounds. See 13.3(14).

- The recommended level of support for the Address attribute should be followed. See 13.3(19). 26/2
- The recommended level of support for the Alignment attribute should be followed. See 13.3(35). 27/2
- The Size of an array object should not include its bounds. See 13.3(41.1/2). 28/2
- If the Size of a subtype allows for efficient independent addressability, then the Size of most objects of the subtype should equal the Size of the subtype. See 13.3(52). 29/2
- A Size clause on a composite subtype should not affect the internal layout of components. See 13.3(53). 30/2
- The recommended level of support for the Size attribute should be followed. See 13.3(56). 31/2
- The recommended level of support for the Component_Size attribute should be followed. See 13.3(73). 32/2
- The recommended level of support for enumeration_representation_clauses should be followed. See 13.4(10). 33/2
- The recommended level of support for record_representation_clauses should be followed. See 13.5.1(22). 34/2
- If a component is represented using a pointer to the actual data of the component which is contiguous with the rest of the object, then the storage place attributes should reflect the place of the actual data. If a component is allocated discontiguously from the rest of the object, then a warning should be generated upon reference to one of its storage place attributes. See 13.5.2(5). 35/2
- The recommended level of support for the nondefault bit ordering should be followed. See 13.5.3(8). 36/2
- Type System.Address should be a private type. See 13.7(37). 37/2
- Operations in System and its children should reflect the target environment; operations that do not make sense should raise Program_Error. See 13.7.1(16). 38/2
- Since the Size of an array object generally does not include its bounds, the bounds should not be part of the converted data in an instance of Unchecked_Conversion. See 13.9(14/2). 39/2
- There should not be unnecessary run-time checks on the result of an Unchecked_Conversion; the result should be returned by reference when possible. Restrictions on Unchecked_Conversions should be avoided. See 13.9(15). 40/2
- The recommended level of support for Unchecked_Conversion should be followed. See 13.9(17). 41/2
- Any cases in which heap storage is dynamically allocated other than as part of the evaluation of an allocator should be documented. See 13.11(23). 42/2
- A default storage pool for an access-to-constant type should not have overhead to support deallocation of individual objects. See 13.11(24). 43/2
- Usually, a storage pool for an access discriminant or access parameter should be created at the point of an allocator, and be reclaimed when the designated object becomes inaccessible. For other anonymous access types, the pool should be created at the point where the type is elaborated and need not support deallocation of individual objects. See 13.11(25). 44/2
- For a standard storage pool, an instance of Unchecked_Deallocation should actually reclaim the storage. See 13.11.2(17). 45/2
- The recommended level of support for the Stream_Size attribute should be followed. See 13.13.2(1.8/2). 46/2

- 47/2 • If not specified, the value of Stream_Size for an elementary type should be the number of bits that corresponds to the minimum number of stream elements required by the first subtype of the type, rounded up to the nearest factor or multiple of the word size that is also a multiple of the stream element size. See 13.13.2(1.6/2).
- 48/2 • If an implementation provides additional named predefined integer types, then the names should end with “Integer”. If an implementation provides additional named predefined floating point types, then the names should end with “Float”. See A.1(52).
- 49/2 • Implementation-defined operations on Wide_Character, Wide_String, Wide_Wide_Character, and Wide_Wide_String should be child units of Wide_Characters or Wide_Wide_Characters. See A.3.1(7/2).
- 50/2 • Bounded string objects should not be implemented by implicit pointers and dynamic allocation. See A.4.4(106).
- 51/2 • Strings.Hash should be good a hash function, returning a wide spread of values for different string values, and similar strings should rarely return the same value. See A.4.9(12/2).
- 52/2 • Any storage associated with an object of type Generator of the random number packages should be reclaimed on exit from the scope of the object. See A.5.2(46).
- 53/2 • Each value of Initiator passed to Reset for the random number packages should initiate a distinct sequence of random numbers, or, if that is not possible, be at least a rapidly varying function of the initiator value. See A.5.2(47).
- 54/2 • Get_Immediate should be implemented with unbuffered input; input should be available immediately; line-editing should be disabled. See A.10.7(23).
- 55/2 • Package Directories.Information should be provided to retrieve other information about a file. See A.16(124/2).
- 56/2 • Directories.Start_Search and Directories.Search should raise Use_Error for malformed patterns. See A.16(125/2).
- 57/2 • Directories.Rename should be supported at least when both New_Name and Old_Name are simple names and New_Name does not identify an existing external file. See A.16(126/2).
- 58/2 • If the execution environment supports subprocesses, the current environment variables should be used to initialize the environment variables of a subprocess. See A.17(32/2).
- 59/2 • Changes to the environment variables made outside the control of Environment_Variables should be reflected immediately. See A.17(33/2).
- 60/2 • Containers.Hash_Type'Modulus should be at least 2^{**32} . Containers.Count_Type'Last should be at least $2^{**31}-1$. See A.18.1(8/2).
- 61/2 • The worst-case time complexity of Element for Containers.Vector should be $O(\log N)$. See A.18.2(256/2).
- 62/2 • The worst-case time complexity of Append with Count = 1 when N is less than the capacity for Containers.Vector should be $O(\log N)$. See A.18.2(257).
- 63/2 • The worst-case time complexity of Prepend with Count = 1 and Delete_First with Count=1 for Containers.Vectors should be $O(N \log N)$. See A.18.2(258/2).
- 64/2 • The worst-case time complexity of a call on procedure Sort of an instance of Containers.Vectors.Generic_Sorting should be $O(N^{**2})$, and the average time complexity should be better than $O(N^{**2})$. See A.18.2(259/2).
- 65/2 • Containers.Vectors.Generic_Sorting.Sort and Containers.Vectors.Generic_Sorting.Merge should minimize copying of elements. See A.18.2(260/2).

- Containers.Vectors.Move should not copy elements, and should minimize copying of internal data structures. See A.18.2(261/2). 66/2
- If an exception is propagated from a vector operation, no storage should be lost, nor any elements removed from a vector unless specified by the operation. See A.18.2(262/2). 67/2
- The worst-case time complexity of Element, Insert with Count=1, and Delete with Count=1 for Containers.Doubly_Linked_Lists should be $O(\log N)$. See A.18.3(160/2). 68/2
- a call on procedure Sort of an instance of Containers.Doubly_Linked_Lists.Generic_Sorting should have an average time complexity better than $O(N^{**}2)$ and worst case no worse than $O(N^{**}2)$. See A.18.3(161/2). 69/2
- Containers.Doubly_Link_Lists.Move should not copy elements, and should minimize copying of internal data structures. See A.18.3(162/2). 70/2
- If an exception is propagated from a list operation, no storage should be lost, nor any elements removed from a list unless specified by the operation. See A.18.3(163/2). 71/2
- Move for a map should not copy elements, and should minimize copying of internal data structures. See A.18.4(83/2). 72/2
- If an exception is propagated from a map operation, no storage should be lost, nor any elements removed from a map unless specified by the operation. See A.18.4(84/2). 73/2
- The average time complexity of Element, Insert, Include, Replace, Delete, Exclude and Find operations that take a key parameter for Containers.Hashed_Maps should be $O(\log N)$. The average time complexity of the subprograms of Containers.Hashed_Maps that take a cursor parameter should be $O(1)$. See A.18.5(62/2). 74/2
- The worst-case time complexity of Element, Insert, Include, Replace, Delete, Exclude and Find operations that take a key parameter for Containers.Ordered_Maps should be $O((\log N)^{**}2)$ or better. The worst-case time complexity of the subprograms of Containers.Ordered_Maps that take a cursor parameter should be $O(1)$. See A.18.6(95/2). 75/2
- Move for sets should not copy elements, and should minimize copying of internal data structures. See A.18.7(104/2). 76/2
- If an exception is propagated from a set operation, no storage should be lost, nor any elements removed from a set unless specified by the operation. See A.18.7(105/2). 77/2
- The average time complexity of the Insert, Include, Replace, Delete, Exclude and Find operations of Containers.Hashed_Sets that take an element parameter should be $O(\log N)$. The average time complexity of the subprograms of Containers.Hashed_Sets that take a cursor parameter should be $O(1)$. The average time complexity of Containers.Hashed_Sets.-Reserve_Capacity should be $O(N)$. See A.18.8(88/2). 78/2
- The worst-case time complexity of the Insert, Include, Replace, Delete, Exclude and Find operations of Containers.Ordered_Sets that take an element parameter should be $O((\log N)^{**}2)$. The worst-case time complexity of the subprograms of Containers.Ordered_Sets that take a cursor parameter should be $O(1)$. See A.18.9(116/2). 79/2
- Containers.Generic_Array_Sort and Containers.Generic_Constrained_Array_Sort should have an average time complexity better than $O(N^{**}2)$ and worst case no worse than $O(N^{**}2)$. See A.18.16(10/2). 80/2
- Containers.Generic_Array_Sort and Containers.Generic_Constrained_Array_Sort should minimize copying of elements. See A.18.16(11/2). 81/2

- 82/2 • If **pragma Export** is supported for a language, the main program should be able to be written in that language. Subprograms named "adainit" and "adafinal" should be provided for elaboration and finalization of the environment task. See B.1(39).
- 83/2 • Automatic elaboration of preelaborated packages should be provided when **pragma Export** is supported. See B.1(40).
- 84/2 • For each supported convention *L* other than Intrinsic, **pragmas Import** and **Export** should be supported for objects of *L*-compatible types and for subprograms, and **pragma Convention** should be supported for *L*-eligible types and for subprograms. See B.1(41).
- 85/2 • If an interface to C, COBOL, or Fortran is provided, the corresponding package or packages described in Annex B, "Interface to Other Languages" should also be provided. See B.2(13).
- 86/2 • The constants `nul`, `wide_nul`, `char16_nul`, and `char32_nul` in package `Interfaces.C` should have a representation of zero. See B.3(62.1/2).
- 87/2 • If C interfacing is supported, the interface correspondences between Ada and C should be supported. See B.3(71).
- 88/2 • If COBOL interfacing is supported, the interface correspondences between Ada and COBOL should be supported. See B.4(98).
- 89/2 • If Fortran interfacing is supported, the interface correspondences between Ada and Fortran should be supported. See B.5(26).
- 90/2 • The machine code or intrinsics support should allow access to all operations normally available to assembly language programmers for the target environment. See C.1(3).
- 91/2 • Interface to assembler should be supported; the default assembler should be associated with the convention identifier `Assembler`. See C.1(4).
- 92/2 • If an entity is exported to assembly language, then the implementation should allocate it at an addressable location even if not otherwise referenced from the Ada code. A call to a machine code or assembler subprogram should be treated as if it could read or update every object that is specified as exported. See C.1(5).
- 93/2 • Little or no overhead should be associated with calling intrinsic and machine-code subprograms. See C.1(10).
- 94/2 • Intrinsic subprograms should be provided to access any machine operations that provide special capabilities or efficiency not normally available. See C.1(16).
- 95/2 • If the `Ceiling_Locking` policy is not in effect and the target system allows for finer-grained control of interrupt blocking, a means for the application to specify which interrupts are to be blocked during protected actions should be provided. See C.3(28/2).
- 96/2 • Interrupt handlers should be called directly by the hardware. See C.3.1(20).
- 97/2 • Violations of any implementation-defined restrictions on interrupt handlers should be detected before run time. See C.3.1(21).
- 98/2 • If implementation-defined forms of interrupt handler procedures are supported, then for each such form of a handler, a type analogous to `Parameterless_Handler` should be specified in a child package of `Interrupts`, with the same operations as in the predefined package `Interrupts`. See C.3.2(25).
- 99/2 • Preelaborated packages should be implemented such that little or no code is executed at run time for the elaboration of entities. See C.4(14).
- 100/2 • If **pragma Discard_Names** applies to an entity, then the amount of storage used for storing names associated with that entity should be reduced. See C.5(8).

- A load or store of a volatile object whose size is a multiple of System.Storage_Unit and whose alignment is nonzero, should be implemented by accessing exactly the bits of the object and no others. See C.6(22/2). 101/2
- A load or store of an atomic object should be implemented by a single load or store instruction. See C.6(23/2). 102/2
- Finalization of task attributes and reclamation of associated storage should be performed as soon as possible after task termination. See C.7.2(30.1/2). 103/2
- If the target domain requires deterministic memory use at run time, storage for task attributes should be pre-allocated statically and the number of attributes pre-allocated should be documented. See C.7.2(30). 104/2
- Names that end with “_Locking” should be used for implementation-defined locking policies. See D.3(17). 105/2
- Names that end with “_Queuing” should be used for implementation-defined queuing policies. See D.4(16). 106/2
- The `abort_statement` should not require the task executing the statement to block. See D.6(9). 107/2
- On a multi-processor, the delay associated with aborting a task on another processor should be bounded. See D.6(10). 108/2
- When feasible, specified restrictions should be used to produce a more efficient implementation. See D.7(21). 109/2
- When appropriate, mechanisms to change the value of Tick should be provided. See D.8(47). 110/2
- Calendar.Clock and Real_Time.Clock should be transformations of the same time base. See D.8(48). 111/2
- The “best” time base which exists in the underlying system should be available to the application through Real_Time.Clock. See D.8(49). 112/2
- When appropriate, implementations should provide configuration mechanisms to change the value of Execution_Time.CPU_Tick. See D.14(29/2). 113/2
- For a timing event, the handler should be executed directly by the real-time clock interrupt mechanism. See D.15(25). 114/2
- The PCS should allow for multiple tasks to call the RPC-receiver. See E.5(28). 115/2
- The System.RPC.Write operation should raise Storage_Error if it runs out of space when writing an item. See E.5(29). 116/2
- If COBOL (respectively, C) is supported in the target environment, then interfacing to COBOL (respectively, C) should be supported as specified in Annex B. See F(7). 117/2
- Packed decimal should be used as the internal representation for objects of subtype *S* when *S*'Machine_Radix = 10. See F.1(2). 118/2
- If Fortran (respectively, C) is supported in the target environment, then interfacing to Fortran (respectively, C) should be supported as specified in Annex B. See G(7). 119/2
- Mixed real and complex operations (as well as pure-imaginary and complex operations) should not be performed by converting the real (resp. pure-imaginary) operand to complex. See G.1.1(56). 120/2
- If RealSigned_Zeros is true for Numerics.Generic_Complex_Types, a rational treatment of the signs of zero results and result components should be provided. See G.1.1(58). 121/2

- 122/2 • If Complex_Types.Real'Signed_Zeros is true for Numerics.Generic_Complex_Elementary_- Functions, a rational treatment of the signs of zero results and result components should be provided. See G.1.2(49).
- 123/2 • For elementary functions, the forward trigonometric functions without a Cycle parameter should not be implemented by calling the corresponding version with a Cycle parameter. Log without a Base parameter should not be implemented by calling Log with a Base parameter. See G.2.4(19).
- 124/2 • For complex arithmetic, the Compose_From_Polar function without a Cycle parameter should not be implemented by calling Compose_From_Polar with a Cycle parameter. See G.2.6(15).
- 125/2 • Solve and Inverse for Numerics.Generic_Real_Arrays should be implemented using established techniques such as LU decomposition and the result should be refined by an iteration on the residuals. See G.3.1(88/2).
- 126/2 • The equality operator should be used to test that a matrix in Numerics.Generic_Real_Matrix is symmetric. See G.3.1(90/2).
- 127/2 • Solve and Inverse for Numerics.Generic_Complex_Arrays should be implemented using established techniques and the result should be refined by an iteration on the residuals. See G.3.2(158/2).
- 128/2 • The equality and negation operators should be used to test that a matrix is Hermitian. See G.3.2(160/2).
- 129/2 • Mixed real and complex operations should not be performed by converting the real operand to complex. See G.3.2(161/2).
- 130/2 • The information produced by pragma Reviewable should be provided in both a human-readable and machine-readable form, and the latter form should be documented. See H.3.1(19).
- 131/2 • Object code listings should be provided both in a symbolic format and in a numeric format. See H.3.1(20).
- 132/2 • If the partition elaboration policy is Sequential and the Environment task becomes permanently blocked during elaboration then the partition should be immediately terminated. See H.6(15/2).

Annex N (informative) Glossary

This Annex contains informal descriptions of some of the terms used in this International Standard. The index provides references to more formal definitions of all of the terms used in this International Standard.	1/2
Abstract type. An abstract type is a tagged type intended for use as an ancestor of other types, but which is not allowed to have objects of its own.	1.1/2
Access type. An access type has values that designate aliased objects. Access types correspond to “pointer types” or “reference types” in some other languages.	2
Aliased. An aliased view of an object is one that can be designated by an access value. Objects allocated by allocators are aliased. Objects can also be explicitly declared as aliased with the reserved word aliased . The Access attribute can be used to create an access value designating an aliased object.	3
Ancestor. An ancestor of a type is the type itself or, in the case of a type derived from other types, its parent type or one of its progenitor types or one of their ancestors. Note that ancestor and descendant are inverse relationships.	3.1/2
Array type. An array type is a composite type whose components are all of the same type. Components are selected by indexing.	4
Category (of types). A category of types is a set of types with one or more common properties, such as primitive operations. A category of types that is closed under derivation is also known as a <i>class</i> .	4.1/2
Character type. A character type is an enumeration type whose values include characters.	5
Class (of types). A class is a set of types that is closed under derivation, which means that if a given type is in the class, then all types derived from that type are also in the class. The set of types of a class share common properties, such as their primitive operations.	6/2
Compilation unit. The text of a program can be submitted to the compiler in one or more compilations. Each compilation is a succession of compilation_units. A compilation_unit contains either the declaration, the body, or a renaming of a program unit.	7
Composite type. A composite type may have components.	8/2
Construct. A <i>construct</i> is a piece of text (explicit or implicit) that is an instance of a syntactic category defined under “Syntax”.	9
Controlled type. A controlled type supports user-defined assignment and finalization. Objects are always finalized before being destroyed.	10
Declaration. A <i>declaration</i> is a language construct that associates a name with (a view of) an entity. A declaration may appear explicitly in the program text (an <i>explicit</i> declaration), or may be supposed to occur at a given place in the text as a consequence of the semantics of another construct (an <i>implicit</i> declaration).	11
<i>This paragraph was deleted.</i>	12/2
Derived type. A derived type is a type defined in terms of one or more other types given in a derived type definition. The first of those types is the parent type of the derived type and any others are progenitor	13/2

types. Each class containing the parent type or a progenitor type also contains the derived type. The derived type inherits properties such as components and primitive operations from the parent and progenitors. A type together with the types derived from it (directly or indirectly) form a derivation class.

13.1/2 **Descendant.** A type is a descendant of itself, its parent and progenitor types, and their ancestors. Note that descendant and ancestor are inverse relationships.

14 **Discrete type.** A discrete type is either an integer type or an enumeration type. Discrete types may be used, for example, in `case_statements` and as array indices.

15/2 **Discriminant.** A discriminant is a parameter for a composite type. It can control, for example, the bounds of a component of the type if the component is an array. A discriminant for a task type can be used to pass data to a task of the type upon creation.

15.1/2 **Elaboration.** The process by which a declaration achieves its run-time effect is called elaboration. Elaboration is one of the forms of execution.

16 **Elementary type.** An elementary type does not have components.

17 **Enumeration type.** An enumeration type is defined by an enumeration of its values, which may be named by identifiers or character literals.

17.1/2 **Evaluation.** The process by which an expression achieves its run-time effect is called evaluation. Evaluation is one of the forms of execution.

18 **Exception.** An *exception* represents a kind of exceptional situation; an occurrence of such a situation (at run time) is called an *exception occurrence*. To *raise* an exception is to abandon normal program execution so as to draw attention to the fact that the corresponding situation has arisen. Performing some actions in response to the arising of an exception is called *handling* the exception.

19 **Execution.** The process by which a construct achieves its run-time effect is called *execution*. Execution of a declaration is also called *elaboration*. Execution of an expression is also called *evaluation*.

19.1/2 **Function.** A function is a form of subprogram that returns a result and can be called as part of an expression.

20 **Generic unit.** A generic unit is a template for a (nongeneric) program unit; the template can be parameterized by objects, types, subprograms, and packages. An instance of a generic unit is created by a `generic_instantiation`. The rules of the language are enforced when a generic unit is compiled, using a generic contract model; additional checks are performed upon instantiation to verify the contract is met. That is, the declaration of a generic unit represents a contract between the body of the generic and instances of the generic. Generic units can be used to perform the role that macros sometimes play in other languages.

20.1/2 **Incomplete type.** An incomplete type gives a view of a type that reveals only some of its properties. The remaining properties are provided by the full view given elsewhere. Incomplete types can be used for defining recursive data structures.

21 **Integer type.** Integer types comprise the signed integer types and the modular types. A signed integer type has a base range that includes both positive and negative numbers, and has operations that may raise an exception when the result is outside the base range. A modular type has a base range whose lower bound is zero, and has operations with “wraparound” semantics. Modular types subsume what are called “unsigned types” in some other languages.

Interface type. An interface type is a form of abstract tagged type which has no components or concrete operations except possibly null procedures. Interface types are used for composing other interfaces and tagged types and thereby provide multiple inheritance. Only an interface type can be used as a progenitor of another type.	21.1/2
Library unit. A library unit is a separately compiled program unit, and is always a package, subprogram, or generic unit. Library units may have other (logically nested) library units as children, and may have other program units physically nested within them. A root library unit, together with its children and grandchildren and so on, form a <i>subsystem</i> .	22
Limited type. A limited type is a type for which copying (such as in an <code>assignment_statement</code>) is not allowed. A nonlimited type is a type for which copying is allowed.	23/2
Object. An object is either a constant or a variable. An object contains a value. An object is created by an <code>object_declaration</code> or by an <code>allocator</code> . A formal parameter is (a view of) an object. A subcomponent of an object is an object.	24
Overriding operation. An overriding operation is one that replaces an inherited primitive operation. Operations may be marked explicitly as overriding or not overriding.	24.1/2
Package. Packages are program units that allow the specification of groups of logically related entities. Typically, a package contains the declaration of a type (often a private type or private extension) along with the declarations of primitive subprograms of the type, which can be called from outside the package, while their inner workings remain hidden from outside users.	25
Parent. The parent of a derived type is the first type given in the definition of the derived type. The parent can be almost any kind of type, including an interface type.	25.1/2
Partition. A <i>partition</i> is a part of a program. Each partition consists of a set of library units. Each partition may run in a separate address space, possibly on a separate computer. A program may contain just one partition. A distributed program typically contains multiple partitions, which can execute concurrently.	26
Pragma. A pragma is a compiler directive. There are language-defined pragmas that give instructions for optimization, listing control, etc. An implementation may support additional (implementation-defined) pragmas.	27
Primitive operations. The primitive operations of a type are the operations (such as subprograms) declared together with the type declaration. They are inherited by other types in the same class of types. For a tagged type, the primitive subprograms are dispatching subprograms, providing run-time polymorphism. A dispatching subprogram may be called with statically tagged operands, in which case the subprogram body invoked is determined at compile time. Alternatively, a dispatching subprogram may be called using a dispatching call, in which case the subprogram body invoked is determined at run time.	28
Private extension. A private extension is a type that extends another type, with the additional properties hidden from its clients.	29/2
Private type. A private type gives a view of a type that reveals only some of its properties. The remaining properties are provided by the full view given elsewhere. Private types can be used for defining abstractions that hide unnecessary details from their clients.	30/2
Procedure. A procedure is a form of subprogram that does not return a result and can only be called by a statement.	30.1/2

- 30.2/2 **Progenitor.** A progenitor of a derived type is one of the types given in the definition of the derived type other than the first. A progenitor is always an interface type. Interfaces, tasks, and protected types may also have progenitors.
- 31 **Program.** A *program* is a set of *partitions*, each of which may execute in a separate address space, possibly on a separate computer. A partition consists of a set of library units.
- 32 **Program unit.** A *program unit* is either a package, a task unit, a protected unit, a protected entry, a generic unit, or an explicitly declared subprogram other than an enumeration literal. Certain kinds of program units can be separately compiled. Alternatively, they can appear physically nested within other program units.
- 33/2 **Protected type.** A protected type is a composite type whose components are accessible only through one of its protected operations which synchronize concurrent access by multiple tasks.
- 34 **Real type.** A real type has values that are approximations of the real numbers. Floating point and fixed point types are real types.
- 35 **Record extension.** A record extension is a type that extends another type by adding additional components.
- 36 **Record type.** A record type is a composite type consisting of zero or more named components, possibly of different types.
- 36.1/2 **Renaming.** A *renaming_declaration* is a declaration that does not define a new entity, but instead defines a view of an existing entity.
- 37 **Scalar type.** A scalar type is either a discrete type or a real type.
- 37.1/2 **Subprogram.** A subprogram is a section of a program that can be executed in various contexts. It is invoked by a subprogram call that may qualify the effect of the subprogram through the passing of parameters. There are two forms of subprograms: functions, which return values, and procedures, which do not.
- 38/2 **Subtype.** A subtype is a type together with a constraint or null exclusion, which constrains the values of the subtype to satisfy a certain condition. The values of a subtype are a subset of the values of its type.
- 38.1/2 **Synchronized.** A synchronized entity is one that will work safely with multiple tasks at one time. A synchronized interface can be an ancestor of a task or a protected type. Such a task or protected type is called a synchronized tagged type.
- 39 **Tagged type.** The objects of a tagged type have a run-time type tag, which indicates the specific type with which the object was originally created. An operand of a class-wide tagged type can be used in a dispatching call; the tag indicates which subprogram body to invoke. Nondispatching calls, in which the subprogram body to invoke is determined at compile time, are also allowed. Tagged types may be extended with additional components.
- 40/2 **Task type.** A task type is a composite type used to represent active entities which execute concurrently and which can communicate via queued task entries. The top-level task of a partition is called the environment task.
- 41/2 **Type.** Each object has a type. A *type* has an associated set of values, and a set of *primitive operations* which implement the fundamental aspects of its semantics. Types are grouped into *categories*. Most language-defined categories of types are also *classes* of types.

View. A view of an entity reveals some or all of the properties of the entity. A single entity may have 42/2 multiple views.

Annex P (informative) Syntax Summary

This Annex summarizes the complete syntax of the language. See 1.1.4 for a description of the notation used.

```

2.3:
identifier ::= 
    identifier_start {identifier_start | identifier_extend}

2.3:
identifier_start ::= 
    letter_uppercase
| letter_lowercase
| letter_titlecase
| letter_modifier
| letter_other
| number_letter

2.3:
identifier_extend ::= 
    mark_non_spacing
| mark_spacing_combining
| number_decimal
| punctuation_connector
| other_format

2.4:
numeric_literal ::= decimal_literal | based_literal

2.4.1:
decimal_literal ::= numeral [.numeral] [exponent]

2.4.1:
numeral ::= digit {[underline] digit}

2.4.1:
exponent ::= E [+] numeral | E – numeral

2.4.1:
digit ::= 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9

2.4.2:
based_literal ::= 
    base # based_numeral [.based_numeral] # [exponent]

2.4.2:
base ::= numeral

2.4.2:
based_numeral ::= 
    extended_digit {[underline] extended_digit}

2.4.2:
extended_digit ::= digit | A | B | C | D | E | F

2.5:
character_literal ::= 'graphic_character'

2.6:
string_literal ::= "{string_element}"

2.6:
string_element ::= "" | non_quotation_mark_graphic_character

2.7:
comment ::= --{non_end_of_line_character}

```

```

2.8:
pragma ::=

  pragma identifier [(pragma_argument_association {, pragma_argument_association})];

2.8:
pragma_argument_association ::=

  [pragma_argument_identifier =>] name
  | [pragma_argument_identifier =>] expression

3.1:
basic_declaration ::=

  type_declaration      | subtype_declaration
  | object_declaration   | number_declaration
  | subprogram_declaration | abstract_subprogram_declaration
  | null_procedure_declaration | package_declaration
  | renaming_declaration | exception_declaration
  | generic_declaration  | generic_instantiation

3.1:
defining_identifier ::= identifier

3.2.1:
type_declaration ::= full_type_declaration
  | incomplete_type_declaration
  | private_type_declaration
  | private_extension_declaration

3.2.1:
full_type_declaration ::=

  type defining_identifier [known_discriminant_part] is type_definition;
  | task_type_declaration
  | protected_type_declaration

3.2.1:
type_definition ::=

  enumeration_type_definition  | integer_type_definition
  | real_type_definition       | array_type_definition
  | record_type_definition     | access_type_definition
  | derived_type_definition    | interface_type_definition

3.2.2:
subtype_declaration ::=

  subtype defining_identifier is subtype_indication;

3.2.2:
subtype_indication ::= [null_exclusion] subtype_mark [constraint]

3.2.2:
subtype_mark ::= subtype_name

3.2.2:
constraint ::= scalar_constraint | composite_constraint

3.2.2:
scalar_constraint ::=

  range_constraint | digits_constraint | delta_constraint

3.2.2:
composite_constraint ::=

  index_constraint | discriminant_constraint

3.3.1:
object_declaration ::=

  defining_identifier_list : [aliased] [constant] subtype_indication [= expression];
  | defining_identifier_list : [aliased] [constant] access_definition [= expression];
  | defining_identifier_list : [aliased] [constant] array_type_definition [= expression];
  | single_task_declaration
  | single_protected_declaration

3.3.1:
defining_identifier_list ::=

  defining_identifier {, defining_identifier}

```

```

3.3.2:
number_declaration ::= 
  defining_identifier_list : constant := static_expression;

3.4:
derived_type_definition ::= 
  [abstract] [limited] new parent_subtype_indication [[and interface_list] record_extension_part]

3.5:
range_constraint ::= range range

3.5:
range ::= range_attribute_reference
  | simple_expression .. simple_expression

3.5.1:
enumeration_type_definition ::= 
  (enumeration_literal_specification {, enumeration_literal_specification})

3.5.1:
enumeration_literal_specification ::= defining_identifier | defining_character_literal

3.5.1:
defining_character_literal ::= character_literal

3.5.4:
integer_type_definition ::= signed_integer_type_definition | modular_type_definition

3.5.4:
signed_integer_type_definition ::= range static_simple_expression .. static_simple_expression

3.5.4:
modular_type_definition ::= mod static_expression

3.5.6:
real_type_definition ::= 
  floating_point_definition | fixed_point_definition

3.5.7:
floating_point_definition ::= 
  digits static_expression [real_range_specification]

3.5.7:
real_range_specification ::= 
  range static_simple_expression .. static_simple_expression

3.5.9:
fixed_point_definition ::= ordinary_fixed_point_definition | decimal_fixed_point_definition

3.5.9:
ordinary_fixed_point_definition ::= 
  delta static_expression real_range_specification

3.5.9:
decimal_fixed_point_definition ::= 
  delta static_expression digits static_expression [real_range_specification]

3.5.9:
digits_constraint ::= 
  digits static_expression [range_constraint]

3.6:
array_type_definition ::= 
  unconstrained_array_definition | constrained_array_definition

3.6:
unconstrained_array_definition ::= 
  array(index_subtype_definition {, index_subtype_definition}) of component_definition

3.6:
index_subtype_definition ::= subtype_mark range <>

3.6:
constrained_array_definition ::= 
  array(discrete_subtype_definition {, discrete_subtype_definition}) of component_definition

```

```

3.6:
discrete_subtype_definition ::= discrete_subtype_indication | range

3.6:
component_definition ::=  

  [aliased] subtype_indication  

  | [aliased] access_definition

3.6.1:
index_constraint ::= (discrete_range {, discrete_range})

3.6.1:
discrete_range ::= discrete_subtype_indication | range

3.7:
discriminant_part ::= unknown_discriminant_part | known_discriminant_part

3.7:
unknown_discriminant_part ::= (<>)

3.7:
known_discriminant_part ::=  

  (discriminant_specification {; discriminant_specification})

3.7:
discriminant_specification ::=  

  defining_identifier_list : [null_exclusion] subtype_mark [:= default_expression]  

  | defining_identifier_list : access_definition [:= default_expression]

3.7:
default_expression ::= expression

3.7.1:
discriminant_constraint ::=  

  (discriminant_association {, discriminant_association})

3.7.1:
discriminant_association ::=  

  [discriminant_selector_name {; discriminant_selector_name} =>] expression

3.8:
record_type_definition ::= [[abstract] tagged] [limited] record_definition

3.8:
record_definition ::=  

  record  

    component_list  

  end record  

  | null record

3.8:
component_list ::=  

  component_item {component_item}  

  | {component_item} variant_part  

  | null;

3.8:
component_item ::= component_declaration | aspect_clause

3.8:
component_declaration ::=  

  defining_identifier_list : component_definition [:= default_expression];

3.8.1:
variant_part ::=  

  case discriminant_direct_name is  

    variant  

    {variant}  

  end case;

```

```

3.8.1:
variant ::= 
  when discrete_choice_list =>
    component_list

3.8.1:
discrete_choice_list ::= discrete_choice { discrete_choice }

3.8.1:
discrete_choice ::= expression | discrete_range | others

3.9.1:
record_extension_part ::= with record_definition

3.9.3:
abstract_subprogram_declaration ::= 
  [overriding_indicator]
  subprogram_specification is abstract;

3.9.4:
interface_type_definition ::= 
  [limited | task | protected | synchronized] interface [and interface_list]

3.9.4:
interface_list ::= interface_subtype_mark {and interface_subtype_mark}

3.10:
access_type_definition ::= 
  [null_exclusion] access_to_object_definition
  | [null_exclusion] access_to_subprogram_definition

3.10:
access_to_object_definition ::= 
  access [general_access_modifier] subtype_indication

3.10:
general_access_modifier ::= all | constant

3.10:
access_to_subprogram_definition ::= 
  access [protected] procedure parameter_profile
  | access [protected] function parameter_and_result_profile

3.10:
null_exclusion ::= not null

3.10:
access_definition ::= 
  [null_exclusion] access [constant] subtype_mark
  | [null_exclusion] access [protected] procedure parameter_profile
  | [null_exclusion] access [protected] function parameter_and_result_profile

3.10.1:
incomplete_type_declaration ::= type defining_identifier [discriminant_part] [is tagged];

3.11:
declarative_part ::= {declarative_item}

3.11:
declarative_item ::= 
  basic_declarative_item | body

3.11:
basic_declarative_item ::= 
  basic_declaration | aspect_clause | use_clause

3.11:
body ::= proper_body | body_stub

3.11:
proper_body ::= 
  subprogram_body | package_body | task_body | protected_body

```

```

4.1:
name ::= direct_name | explicit_dereference
| indexed_component | slice
| selected_component | attribute_reference
| type_conversion | function_call
| character_literal

4.1:
direct_name ::= identifier | operator_symbol

4.1:
prefix ::= name | implicit_dereference

4.1:
explicit_dereference ::= name.all

4.1:
implicit_dereference ::= name

4.1.1:
indexed_component ::= prefix(expression {, expression})

4.1.2:
slice ::= prefix(discrete_range)

4.1.3:
selected_component ::= prefix . selector_name

4.1.3:
selector_name ::= identifier | character_literal | operator_symbol

4.1.4:
attribute_reference ::= prefix'attribute_designator

4.1.4:
attribute_designator ::= identifier[(static_expression)]
| Access | Delta | Digits

4.1.4:
range_attribute_reference ::= prefix'range_attribute_designator

4.1.4:
range_attribute_designator ::= Range[(static_expression)]

4.3:
aggregate ::= record_aggregate | extension_aggregate | array_aggregate

4.3.1:
record_aggregate ::= (record_component_association_list)

4.3.1:
record_component_association_list ::= record_component_association {, record_component_association}
| null record

4.3.1:
record_component_association ::= [component_choice_list =>] expression
| component_choice_list => <>

4.3.1:
component_choice_list ::= component_selector_name {& component_selector_name}
| others

4.3.2:
extension_aggregate ::= (ancestor_part with record_component_association_list)

4.3.2:
ancestor_part ::= expression | subtype_mark

```

```

4.3.3:
array_aggregate ::= positional_array_aggregate | named_array_aggregate
4.3.3:
positional_array_aggregate ::= (expression, expression {, expression})
| (expression {, expression}, others => expression)
| (expression {, expression}, others => <>)
4.3.3:
named_array_aggregate ::= (array_component_association {, array_component_association})
4.3.3:
array_component_association ::= discrete_choice_list => expression
| discrete_choice_list => <>
4.4:
expression ::= relation {and relation} | relation {and then relation}
| relation {or relation} | relation {or else relation}
| relation {xor relation}
4.4:
relation ::= simple_expression [relational_operator simple_expression]
| simple_expression [not] in range
| simple_expression [not] in subtype_mark
4.4:
simple_expression ::= [unary_adding_operator] term {binary_adding_operator term}
4.4:
term ::= factor {multiplying_operator factor}
4.4:
factor ::= primary [** primary] | abs primary | not primary
4.4:
primary ::= numeric_literal | null | string_literal | aggregate
| name | qualified_expression | allocator | (expression)
4.5:
logical_operator ::= and | or | xor
4.5:
relational_operator ::= = | /= | < | <= | > | >=
4.5:
binary_adding_operator ::= + | - | &
amp;
4.5:
unary_adding_operator ::= + | -
4.5:
multiplying_operator ::= * | / | mod | rem
4.5:
highest_precedence_operator ::= ** | abs | not
4.6:
type_conversion ::= subtype_mark(expression)
| subtype_mark(name)
4.7:
qualified_expression ::= subtype_mark'(expression) | subtype_mark'aggregate

```

```

4.8:
allocator ::= new subtype_indication | new qualified_expression

5.1:
sequence_of_statements ::= statement {statement}

5.1:
statement ::= {label} simple_statement | {label} compound_statement

5.1:
simple_statement ::= null_statement
| assignment_statement      | exit_statement
| goto_statement            | procedure_call_statement
| simple_return_statement   | entry_call_statement
| requeue_statement         | delay_statement
| abort_statement           | raise_statement
| code_statement

5.1:
compound_statement ::= if_statement          | case_statement
| loop_statement             | block_statement
| extended_return_statement | accept_statement
| accept_statement          | select_statement

5.1:
null_statement ::= null;

5.1:
label ::= <<label_statement_identifier>>

5.1:
statement_identifier ::= direct_name

5.2:
assignment_statement ::= variable_name := expression;

5.3:
if_statement ::= if condition then
  sequence_of_statements
{elsif condition then
  sequence_of_statements}
[else
  sequence_of_statements]
end if;

5.3:
condition ::= boolean_expression

5.4:
case_statement ::= case expression is
  case_statement_alternative
  {case_statement_alternative}
end case;

5.4:
case_statement_alternative ::= when discrete_choice_list =>
  sequence_of_statements

5.5:
loop_statement ::= [loop_statement_identifier:] [iteration_scheme] loop
  sequence_of_statements
end loop [loop_identifier];

```

```

5.5:
iteration_scheme ::= while condition
| for loop_parameter_specification

5.5:
loop_parameter_specification ::= 
  defining_identifier in [reverse] discrete_subtype_definition

5.6:
block_statement ::= 
  [block_statement_identifier]
  [declare
   declarative_part]
  begin
   handled_sequence_of_statements
  end [block_identifier];

5.7:
exit_statement ::= 
  exit [loop_name] [when condition];

5.8:
goto_statement ::= goto label_name;

6.1:
subprogram_declaration ::= 
  [overriding_indicator]
  subprogram_specification;

6.1:
subprogram_specification ::= 
  procedure_specification
| function_specification

6.1:
procedure_specification ::= procedure defining_program_unit_name parameter_profile

6.1:
function_specification ::= function defining_designator parameter_and_result_profile

6.1:
designator ::= [parent_unit_name . ]identifier | operator_symbol

6.1:
defining_designator ::= defining_program_unit_name | defining_operator_symbol

6.1:
defining_program_unit_name ::= [parent_unit_name . ]defining_identifier

6.1:
operator_symbol ::= string_literal

6.1:
defining_operator_symbol ::= operator_symbol

6.1:
parameter_profile ::= [formal_part]

6.1:
parameter_and_result_profile ::= 
  [formal_part] return [null_exclusion] subtype_mark
| [formal_part] return access_definition

6.1:
formal_part ::= 
  (parameter_specification {; parameter_specification})

6.1:
parameter_specification ::= 
  defining_identifier_list : mode [null_exclusion] subtype_mark [= default_expression]
| defining_identifier_list : access_definition [= default_expression]

6.1:
mode ::= [in] | in out | out

```

```

6.3:
subprogram_body ::= 
  [overriding_indicator]
  subprogram_specification is
    declarative_part
  begin
    handled_sequence_of_statements
  end [designator];

6.4:
procedure_call_statement ::= 
  procedure_name;
| procedure_prefix actual_parameter_part;

6.4:
function_call ::= 
  function_name
| function_prefix actual_parameter_part

6.4:
actual_parameter_part ::= 
  (parameter_association {, parameter_association} )

6.4:
parameter_association ::= 
  [formal_parameter_selector_name =>] explicit_actual_parameter

6.4:
explicit_actual_parameter ::= expression | variable_name

6.5:
simple_return_statement ::= return [expression];

6.5:
extended_return_statement ::= 
  return defining_identifier : [aliased] return_subtype_indication [= expression] [do
    handled_sequence_of_statements
  end return];

6.5:
return_subtype_indication ::= subtype_indication | access_definition

6.7:
null_procedure_declaration ::= 
  [overriding_indicator]
  procedure_specification is null;

7.1:
package_declaration ::= package_specification;

7.1:
package_specification ::= 
  package defining_program_unit_name is
    {basic_declarative_item}
  [private
    {basic_declarative_item}]
  end [[parent_unit_name.]identifier]

7.2:
package_body ::= 
  package body defining_program_unit_name is
    declarative_part
  [begin
    handled_sequence_of_statements]
  end [[parent_unit_name.]identifier];

7.3:
private_type_declaration ::= 
  type defining_identifier [discriminant_part] is [[abstract] tagged] [limited] private;

```

```

7.3:
private_extension_declaration ::=

  type defining_identifier [discriminant_part] is
    [abstract] [limited | synchronized] new ancestor_subtype_indication
    [and interface_list] with private;

8.3.1:
overriding_indicator ::= [not] overriding

8.4:
use_clause ::= use_package_clause | use_type_clause

8.4:
use_package_clause ::= use package_name {, package_name};

8.4:
use_type_clause ::= use type subtype_mark {, subtype_mark};

8.5:
renaming_declaration ::=
  object_renaming_declaration
  | exception_renaming_declaration
  | package_renaming_declaration
  | subprogram_renaming_declaration
  | generic_renaming_declaration

8.5.1:
object_renaming_declaration ::=

  defining_identifier : [null_exclusion] subtype_mark renames object_name;
  | defining_identifier : access_definition renames object_name;

8.5.2:
exception_renaming_declaration ::= defining_identifier : exception renames exception_name;

8.5.3:
package_renaming_declaration ::= package defining_program_unit_name renames package_name;

8.5.4:
subprogram_renaming_declaration ::=
  [overriding_indicator]
  subprogram_specification renames callable_entity_name;

8.5.5:
generic_renaming_declaration ::=
  generic package      defining_program_unit_name renames generic_package_name;
  | generic procedure   defining_program_unit_name renames generic_procedure_name;
  | generic function     defining_program_unit_name renames generic_function_name;

9.1:
task_type_declaration ::=
  task type defining_identifier [known_discriminant_part] [is
    [new interface_list with]
    task_definition];

9.1:
single_task_declaration ::=

  task defining_identifier [is
    [new interface_list with]
    task_definition];

9.1:
task_definition ::=

  {task_item}
  [ private
    {task_item}]
  end [task_identifier]

9.1:
task_item ::= entry_declaration | aspect_clause

```

```

9.1:
task_body ::= 
  task body defining_identifier is
    declarative_part
  begin
    handled_sequence_of_statements
  end [task_identifier];

9.4:
protected_type_declaration ::= 
  protected type defining_identifier [known_discriminant_part] is
    [new interface_list with]
    protected_definition;

9.4:
single_protected_declaration ::= 
  protected defining_identifier is
    [new interface_list with]
    protected_definition;

9.4:
protected_definition ::= 
  { protected_operation_declaration }
[ private
  { protected_element_declaration } ]
end [protected_identifier];

9.4:
protected_operation_declaration ::= subprogram_declaration
  | entry_declaration
  | aspect_clause

9.4:
protected_element_declaration ::= protected_operation_declaration
  | component_declaration

9.4:
protected_body ::= 
  protected body defining_identifier is
    { protected_operation_item }
  end [protected_identifier];

9.4:
protected_operation_item ::= subprogram_declaration
  | subprogram_body
  | entry_body
  | aspect_clause

9.5.2:
entry_declaration ::= 
  [overriding_indicator]
  entry defining_identifier [(discrete_subtype_definition)] parameter_profile;

9.5.2:
accept_statement ::= 
  accept entry_direct_name [(entry_index)] parameter_profile [do
    handled_sequence_of_statements
  end [entry_identifier]];

9.5.2:
entry_index ::= expression

9.5.2:
entry_body ::= 
  entry defining_identifier entry_body_formal_part entry_barrier is
    declarative_part
  begin
    handled_sequence_of_statements
  end [entry_identifier];

```

```

9.5.2:
entry_body_formal_part ::= [(entry_index_specification)] parameter_profile
9.5.2:
entry_barrier ::= when condition
9.5.2:
entry_index_specification ::= for defining_identifier in discrete_subtype_definition
9.5.3:
entry_call_statement ::= entry_name [actual_parameter_part];
9.5.4:
enqueue_statement ::= enqueue entry_name [with abort];
9.6:
delay_statement ::= delay_until_statement | delay_relative_statement
9.6:
delay_until_statement ::= delay until delay_expression;
9.6:
delay_relative_statement ::= delay delay_expression;
9.7:
select_statement ::= 
  selective_accept
  | timed_entry_call
  | conditional_entry_call
  | asynchronous_select
9.7.1:
selective_accept ::= 
  select
  [guard]
  select_alternative
{ or
  [guard]
  select_alternative }
[ else
  sequence_of_statements ]
end select;
9.7.1:
guard ::= when condition =>
9.7.1:
select_alternative ::= 
  accept_alternative
  | delay_alternative
  | terminate_alternative
9.7.1:
accept_alternative ::= 
  accept_statement [sequence_of_statements]
9.7.1:
delay_alternative ::= 
  delay_statement [sequence_of_statements]
9.7.1:
terminate_alternative ::= terminate;
9.7.2:
timed_entry_call ::= 
  select
  entry_call_alternative
  or
  delay_alternative
end select;

```

9.7.2:
entry_call_alternative ::=
procedure_or_entry_call [sequence_of_statements]

9.7.2:
procedure_or_entry_call ::=
procedure_call_statement | entry_call_statement

9.7.3:
conditional_entry_call ::=
select
entry_call_alternative
else
sequence_of_statements
end select;

9.7.4:
asynchronous_select ::=
select
triggering_alternative
then abort
abortable_part
end select;

9.7.4:
triggering_alternative ::= triggering_statement [sequence_of_statements]

9.7.4:
triggering_statement ::= procedure_or_entry_call | delay_statement

9.7.4:
abortable_part ::= sequence_of_statements

9.8:
abort_statement ::= abort task_name {, task_name};

10.1.1:
compilation ::= {compilation_unit}

10.1.1:
compilation_unit ::=
context_clause library_item
| context_clause subunit

10.1.1:
library_item ::= [private] library_unit_declaration
| library_unit_body
| [private] library_unit_renaming_declaration

10.1.1:
library_unit_declaration ::=
subprogram_declaration | package_declaration
| generic_declaration | generic_instantiation

10.1.1:
library_unit_renaming_declaration ::=
package_renaming_declaration
| generic_renaming_declaration
| subprogram_renaming_declaration

10.1.1:
library_unit_body ::= subprogram_body | package_body

10.1.1:
parent_unit_name ::= name

10.1.2:
context_clause ::= {context_item}

10.1.2:
context_item ::= with_clause | use_clause

```

10.1.2:
with_clause ::= limited_with_clause | nonlimited_with_clause
10.1.2:
limited_with_clause ::= limited [private] with library_unit_name {, library_unit_name};
10.1.2:
nonlimited_with_clause ::= [private] with library_unit_name {, library_unit_name};
10.1.3:
body_stub ::= subprogram_body_stub | package_body_stub | task_body_stub | protected_body_stub
10.1.3:
subprogram_body_stub ::= 
  [overriding_indicator]
  subprogram_specification is separate;
10.1.3:
package_body_stub ::= package body defining_identifier is separate;
10.1.3:
task_body_stub ::= task body defining_identifier is separate;
10.1.3:
protected_body_stub ::= protected body defining_identifier is separate;
10.1.3:
subunit ::= separate (parent_unit_name) proper_body
11.1:
exception_declaration ::= defining_identifier_list : exception;
11.2:
handled_sequence_of_statements ::= 
  sequence_of_statements
  [exception
   exception_handler
   {exception_handler}]
11.2:
exception_handler ::= 
  when [choice_parameter_specification:] exception_choice { exception_choice } =>
  sequence_of_statements
11.2:
choice_parameter_specification ::= defining_identifier
11.2:
exception_choice ::= exception_name | others
11.3:
raise_statement ::= raise;
  | raise exception_name [with string_expression];
12.1:
generic_declaration ::= generic_subprogram_declaration | generic_package_declaration
12.1:
generic_subprogram_declaration ::= 
  generic_formal_part subprogram_specification;
12.1:
generic_package_declaration ::= 
  generic_formal_part package_specification;
12.1:
generic_formal_part ::= generic {generic_formal_parameter_declaration | use_clause};
12.1:
generic_formal_parameter_declaration ::= 
  formal_object_declaration
  | formal_type_declaration
  | formal_subprogram_declaration
  | formal_package_declaration

```

12.3:

generic_instantiation ::=

- package** defining_program_unit_name **is**
- new generic_package_name [generic_actual_part];
- | [overriding_indicator]
- procedure** defining_program_unit_name **is**
- new generic_procedure_name [generic_actual_part];
- | [overriding_indicator]
- function** defining_designator **is**
- new generic_function_name [generic_actual_part];

12.3:

generic_actual_part ::=

- (generic_association {, generic_association})

12.3:

generic_association ::=

- [generic_formal_parameter_selector_name =>] explicit_generic_actual_parameter

12.3:

explicit_generic_actual_parameter ::= expression | variable_name

- | subprogram_name | entry_name | subtype_mark
- | package_instance_name

12.4:

formal_object_declaration ::=

- defining_identifier_list : mode [null_exclusion] subtype_mark [= default_expression];
- defining_identifier_list : mode access_definition [= default_expression];

12.5:

formal_type_declaration ::=

- type** defining_identifier[discriminant_part] **is** formal_type_definition;

12.5:

formal_type_definition ::=

- formal_private_type_definition
- | formal_derived_type_definition
- | formal_discrete_type_definition
- | formal_signed_integer_type_definition
- | formal_modular_type_definition
- | formal_floating_point_definition
- | formal_ordinary_fixed_point_definition
- | formal_decimal_fixed_point_definition
- | formal_array_type_definition
- | formal_access_type_definition
- | formal_interface_type_definition

12.5.1:

formal_private_type_definition ::= [[**abstract**] **tagged**] [**limited**] **private**

12.5.1:

formal_derived_type_definition ::=

- [**abstract**] [**limited** | **synchronized**] new subtype_mark [[**and** interface_list]**with private**]

12.5.2:

formal_discrete_type_definition ::= (<>)

12.5.2:

formal_signed_integer_type_definition ::= **range** <>

12.5.2:

formal_modular_type_definition ::= **mod** <>

12.5.2:

formal_floating_point_definition ::= **digits** <>

12.5.2:

formal_ordinary_fixed_point_definition ::= **delta** <>

12.5.2:

formal_decimal_fixed_point_definition ::= **delta** <> **digits** <>

```

12.5.3:
formal_array_type_definition ::= array_type_definition

12.5.4:
formal_access_type_definition ::= access_type_definition

12.5.5:
formal_interface_type_definition ::= interface_type_definition

12.6:
formal_subprogram_declaration ::= formal_concrete_subprogram_declaration
| formal_abstract_subprogram_declaration

12.6:
formal_concrete_subprogram_declaration ::= 
    with subprogram_specification [is subprogram_default];

12.6:
formal_abstract_subprogram_declaration ::= 
    with subprogram_specification is abstract [subprogram_default];

12.6:
subprogram_default ::= default_name | <> | null

12.6:
default_name ::= name

12.7:
formal_package_declaration ::= 
    with package defining_identifier is new generic_package_name formal_package_actual_part;

12.7:
formal_package_actual_part ::= 
    (others =>) <>
    | [generic_actual_part]
    | (formal_package_association {, formal_package_association} [, others => <>])

12.7:
formal_package_association ::= 
    generic_association
    | generic_formal_parameter_selector_name => <>

13.1:
aspect_clause ::= attribute_definition_clause
    | enumeration_representation_clause
    | record_representation_clause
    | at_clause

13.1:
local_name ::= direct_name
    | direct_name'attribute_designator
    | library_unit_name

13.3:
attribute_definition_clause ::= 
    for local_name'attribute_designator use expression;
    | for local_name'attribute_designator use name;

13.4:
enumeration_representation_clause ::= 
    for first_subtype_local_name use enumeration_aggregate;

13.4:
enumeration_aggregate ::= array_aggregate

13.5.1:
record_representation_clause ::= 
    for first_subtype_local_name use
        record [mod_clause]
        {component_clause}
    end record;

```

```
13.5.1:  
component_clause ::=  
  component_local_name at position range first_bit .. last_bit;  
13.5.1:  
position ::= static_expression  
13.5.1:  
first_bit ::= static_simple_expression  
13.5.1:  
last_bit ::= static_simple_expression  
13.8:  
code_statement ::= qualified_expression;  
13.12:  
restriction ::= restriction_identifier  
  | restriction_parameter_identifier => restriction_parameter_argument  
13.12:  
restriction_parameter_argument ::= name | expression  
J.3:  
delta_constraint ::= delta static_expression [range_constraint]  
J.7:  
at_clause ::= for direct_name use at expression;  
J.8:  
mod_clause ::= at mod static_expression;
```

Syntax Cross Reference

In the following syntax cross reference, each syntactic category is followed by the clause number where it is defined. In addition, each syntactic category *S* is followed by a list of the categories that use *S* in their definitions. For example, the first listing below shows that `abort_statement` appears in the definition of `simple_statement`.

<code>abort_statement</code>	9.8	<code>array_type_definition</code>	3.6
<code>simple_statement</code>	5.1	<code>formal_array_type_definition</code>	12.5.3
<code>abortable_part</code>	9.7.4	<code>object_declaration</code>	3.3.1
<code>asynchronous_select</code>	9.7.4	<code>type_definition</code>	3.2.1
<code>abstract_subprogram_declaration</code>	3.9.3	<code>aspect_clause</code>	13.1
<code>basic_declaration</code>	3.1	<code>basic_declarative_item</code>	3.11
<code>accept_alternative</code>	9.7.1	<code>component_item</code>	3.8
<code>select_alternative</code>	9.7.1	<code>protected_operation_declaration</code>	9.4
<code>accept_statement</code>	9.5.2	<code>protected_operation_item</code>	9.4
<code>accept_alternative</code>	9.7.1	<code>task_item</code>	9.1
<code>compound_statement</code>	5.1	<code>assignment_statement</code>	5.2
<code>access_definition</code>	3.10	<code>simple_statement</code>	5.1
<code>component_definition</code>	3.6	<code>asynchronous_select</code>	9.7.4
<code>discriminant_specification</code>	3.7	<code>select_statement</code>	9.7
<code>formal_object_declaration</code>	12.4	<code>at_clause</code>	J.7
<code>object_declaration</code>	3.3.1	<code>aspect_clause</code>	13.1
<code>object_renaming_declaration</code>	8.5.1	<code>attribute_definition_clause</code>	13.3
<code>parameter_and_result_profile</code>	6.1	<code>aspect_clause</code>	13.1
<code>parameter_specification</code>	6.1	<code>attribute_designator</code>	4.1.4
<code>return_subtype_indication</code>	6.5	<code>attribute_definition_clause</code>	13.3
<code>access_to_object_definition</code>	3.10	<code>attribute_reference</code>	4.1.4
<code>access_type_definition</code>	3.10	<code>local_name</code>	13.1
<code>access_to_subprogram_definition</code>	3.10	<code>attribute_reference</code>	4.1.4
<code>access_type_definition</code>	3.10	<code>name</code>	4.1
<code>access_type_definition</code>	3.10	<code>base</code>	2.4.2
<code>formal_access_type_definition</code>	12.5.4	<code>based_literal</code>	2.4.2
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<code>actual_parameter_part</code>	6.4	<code>numeric_literal</code>	2.4
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type_definition	3.2.1	delay_alternative	9.7.1
record_aggregate	4.3.1	entry_call_alternative	9.7.2
aggregate	4.3	exception_handler	11.2
record_component_association	4.3.1	handled_sequence_of_statements	11.2
record_component_association_list	4.3.1	if_statement	5.3
record_component_association_list	4.3.1	loop_statement	5.5
extension_aggregate	4.3.2	selective_accept	9.7.1
record_aggregate	4.3.1	triggering_alternative	9.7.4
record_definition	3.8	simple_expression	4.4
record_extension_part	3.9.1	first_bit	13.5.1
record_type_definition	3.8	last_bit	13.5.1
record_extension_part	3.9.1	range	3.5
derived_type_definition	3.4	real_range_specification	3.5.7
record_representation_clause	13.5.1	relation	4.4
aspect_clause	13.1	signed_integer_type_definition	3.5.4
record_type_definition	3.8	simple_return_statement	6.5
type_definition	3.2.1	simple_statement	5.1

simple_statement	5.1	subtype_indication	3.2.2
statement	5.1	access_to_object_definition	3.10
single_protected_declaration	9.4	allocator	4.8
object_declaration	3.3.1	component_definition	3.6
single_task_declaration	9.1	derived_type_definition	3.4
object_declaration	3.3.1	discrete_range	3.6.1
slice 4.1.2		discrete_subtype_definition	3.6
name	4.1	object_declaration	3.3.1
statement	5.1	private_extension_declaration	7.3
sequence_of_statements	5.1	return_subtype_indication	6.5
statement_identifier	5.1	subtype_declaration	3.2.2
block_statement	5.6	subtype_mark	3.2.2
label	5.1	access_definition	3.10
loop_statement	5.5	ancestor_part	4.3.2
string_element	2.6	discriminant_specification	3.7
string_literal	2.6	explicit_generic_actual_parameter	12.3
string_literal	2.6	formal_derived_type_definition	12.5.1
operator_symbol	6.1	formal_object_declaration	12.4
primary	4.4	index_subtype_definition	3.6
subprogram_body	6.3	interface_list	3.9.4
library_unit_body	10.1.1	object_renaming_declaration	8.5.1
proper_body	3.11	parameter_and_result_profile	6.1
protected_operation_item	9.4	parameter_specification	6.1
subprogram_body_stub	10.1.3	qualified_expression	4.7
body_stub	10.1.3	relation	4.4
subprogram_declaration	6.1	subtype_indication	3.2.2
basic_declaration	3.1	type_conversion	4.6
library_unit_declaration	10.1.1	use_type_clause	8.4
protected_operation_declaration	9.4	subunit	10.1.3
protected_operation_item	9.4	compilation_unit	10.1.1
subprogram_default	12.6	task_body	9.1
formal_abstract_subprogram_declaration	12.6	proper_body	3.11
formal_concrete_subprogram_declaration	12.6	task_body_stub	10.1.3
subprogram_renaming_declaration	8.5.4	body_stub	10.1.3
library_unit_renaming_declaration	10.1.1	task_definition	9.1
renaming_declaration	8.5	single_task_declaration	9.1
subprogram_specification	6.1	task_type_declaration	9.1
abstract_subprogram_declaration	3.9.3	task_item	9.1
formal_abstract_subprogram_declaration	12.6	task_definition	9.1
formal_concrete_subprogram_declaration	12.6	task_type_declaration	9.1
generic_subprogram_declaration	12.1	term 4.4	4.4
subprogram_body	6.3	simple_expression	4.4
subprogram_body_stub	10.1.3	terminate_alternative	9.7.1
subprogram_declaration	6.1	select_alternative	9.7.1
subprogram_renaming_declaration	8.5.4	timed_entry_call	9.7.2
subtype_declaration	3.2.2	select_statement	9.7
basic_declaration	3.1	triggering_alternative	9.7.4
		asynchronous_select	9.7.4
		triggering_statement	9.7.4
		triggering_alternative	9.7.4
		type_conversion	4.6
		name	4.1
		type_declaration	3.2.1
		basic_declaration	3.1

type_definition	3.2.1	use_clause	8.4
full_type_declaration	3.2.1	basic_declarative_item	3.11
unary_adding_operator	4.5	context_item	10.1.2
simple_expression	4.4	generic_formal_part	12.1
unconstrained_array_definition	3.6	use_package_clause	8.4
array_type_definition	3.6	use_clause	8.4
underline	...	use_type_clause	8.4
based_numeral	2.4.2	use_clause	8.4
numeral	2.4.1	variant	3.8.1
unknown_discriminant_part	3.7	variant_part	3.8.1
discriminant_part	3.7	variant_part	3.8.1
		component_list	3.8
		with_clause	10.1.2
		context_item	10.1.2

Annex Q (informative) Language-Defined Entities

This annex lists the language-defined entities of the language. A list of language-defined library units can be found in Annex A, “Predefined Language Environment”. 1/2

Q.1 Language-Defined Packages

This clause lists all language-defined packages. 1/2

Ada A.2(2)	Containers
Address_To_Access_Conversions	<i>child of</i> Ada A.18.1(3/2)
<i>child of</i> System 13.7.2(2)	Conversions
Arithmetic	<i>child of</i> Ada.Characters A.3.4(2/2)
<i>child of</i> Ada.Calendar 9.6.1(8/2)	Decimal
ASCII	<i>child of</i> Ada F.2(2)
<i>in</i> Standard A.1(36.3/2)	Decimal_Conversions
Assertions	<i>in</i> Interfaces.COBOL B.4(31)
<i>child of</i> Ada 11.4.2(12/2)	Decimal_IO
Asynchronous_Task_Control	<i>in</i> Ada.Text_IO A.10.1(73)
<i>child of</i> Ada D.11(3/2)	Decimal_Output
Bounded	<i>in</i> Ada.Text_IO.Editing F.3.3(11)
<i>child of</i> Ada.Strings A.4.4(3)	Direct_IO
Bounded_IO	<i>child of</i> Ada A.8.4(2)
<i>child of</i> Ada.Text_IO A.10.11(3/2)	Directories
<i>child of</i> Ada.Wide_Text_IO A.11(4/2)	<i>child of</i> Ada A.16(3/2)
<i>child of</i> Ada.Wide_Wide_Text_IO A.11(4/2)	Discrete_Random
C	<i>child of</i> Ada.Numerics A.5.2(17)
<i>child of</i> Interfaces B.3(4)	Dispatching
Calendar	<i>child of</i> Ada D.2.1(1.2/2)
<i>child of</i> Ada 9.6(10)	Doubly_Linked_Lists
Characters	<i>child of</i> Ada.Containers A.18.3(5/2)
<i>child of</i> Ada A.3.1(2)	Dynamic_Priorities
COBOL	<i>child of</i> Ada D.5.1(3/2)
<i>child of</i> Interfaces B.4(7)	EDF
Command_Line	<i>child of</i> Ada.Dispatching D.2.6(9/2)
<i>child of</i> Ada A.15(3)	Editing
Complex_Arrays	<i>child of</i> Ada.Text_IO F.3.3(3)
<i>child of</i> Ada.Numerics G.3.2(53/2)	<i>child of</i> Ada.Wide_Text_IO F.3.4(1)
Complex_Elementary_Functions	<i>child of</i> Ada.Wide_Wide_Text_IO F.3.5(1/2)
<i>child of</i> Ada.Numerics G.1.2(9/1)	Elementary_Functions
Complex_Text_IO	<i>child of</i> Ada.Numerics A.5.1(9/1)
<i>child of</i> Ada G.1.3(9.1/2)	Enumeration_IO
Complex_Types	<i>in</i> Ada.Text_IO A.10.1(79)
<i>child of</i> Ada.Numerics G.1.1(25/1)	Environment_Variables
Complex_IO	<i>child of</i> Ada A.17(3/2)
<i>child of</i> Ada.Text_IO G.1.3(3)	Exceptions
<i>child of</i> Ada.Wide_Text_IO G.1.4(1)	<i>child of</i> Ada 11.4.1(2/2)
<i>child of</i> Ada.Wide_Wide_Text_IO G.1.5(1/2)	Execution_Time
Constants	<i>child of</i> Ada D.14(3/2)
<i>child of</i> Ada.Strings.Maps A.4.6(3/2)	Finalization
	<i>child of</i> Ada 7.6(4/1)

- Fixed
 - child of* Ada.Strings A.4.3(5)
- Fixed_IO
 - in* Ada.Text_IO A.10.1(68)
- Float_Random
 - child of* Ada.Numerics A.5.2(5)
- Float_Text_IO
 - child of* Ada A.10.9(33)
- Float_Wide_Text_IO
 - child of* Ada A.11(2/2)
- Float_Wide_Wide_Text_IO
 - child of* Ada A.11(3/2)
- Float_IO
 - in* Ada.Text_IO A.10.1(63)
- Formatting
 - child of* Ada.Calendar 9.6.1(15/2)
- Fortran
 - child of* Interfaces B.5(4)
- Generic_Complex_Arrays
 - child of* Ada.Numerics G.3.2(2/2)
- Generic_Complex_Elementary_Functions
 - child of* Ada.Numerics G.1.2(2/2)
- Generic_Complex_Types
 - child of* Ada.Numerics G.1.1(2/1)
- Generic_Dispatching_Constructor
 - child of* Ada.Tags 3.9(18.2/2)
- Generic_Elementary_Functions
 - child of* Ada.Numerics A.5.1(3)
- Generic_Bounded_Length
 - in* Ada.Strings.Bounded A.4.4(4)
- Generic_Keys
 - in* Ada.Containers.Hashed_Sets A.18.8(50/2)
 - in* Ada.Containers.Ordered_Sets A.18.9(62/2)
- Generic_Real_Arrays
 - child of* Ada.Numerics G.3.1(2/2)
- Generic_Sorting
 - in* Ada.Containers.Doubly_Linked_Lists A.18.3(47/2)
 - in* Ada.Containers.Vectors A.18.2(75/2)
- Group_Budgets
 - child of* Ada.Execution_Time D.14.2(3/2)
- Handling
 - child of* Ada.Characters A.3.2(2/2)
- Hashed_Maps
 - child of* Ada.Containers A.18.5(2/2)
- Hashed_Sets
 - child of* Ada.Containers A.18.8(2/2)
- Indefinite_Doubly_Linked_Lists
 - child of* Ada.Containers A.18.11(2/2)
- Indefinite_Hashed_Maps
 - child of* Ada.Containers A.18.12(2/2)
- Indefinite_Hashed_Sets
 - child of* Ada.Containers A.18.14(2/2)
- Indefinite_Ordered_Maps
 - child of* Ada.Containers A.18.13(2/2)
- Indefinite_Ordered_Sets
 - child of* Ada.Containers A.18.15(2/2)
- Indefinite_Vectors
 - child of* Ada.Containers A.18.10(2/2)
- Information
 - child of* Ada.Directories A.16(124/2)

- Integer_Text_IO
 - child of* Ada A.10.8(21)
- Integer_Wide_Text_IO
 - child of* Ada A.11(2/2)
- Integer_Wide_Wide_Text_IO
 - child of* Ada A.11(3/2)
- Integer_IO
 - in* Ada.Text_IO A.10.1(52)
- Interfaces B.2(3)
- Interrupts
 - child of* Ada C.3.2(2)
- IO_Exceptions
 - child of* Ada A.13(3)
- Latin_1
 - child of* Ada.Characters A.3.3(3)
- Machine_Code
 - child of* System 13.8(7)
- Maps
 - child of* Ada.Strings A.4.2(3/2)
- Modular_IO
 - in* Ada.Text_IO A.10.1(57)
- Names
 - child of* Ada.Interrupts C.3.2(12)
- Numerics
 - child of* Ada A.5(3/2)
- Ordered_Maps
 - child of* Ada.Containers A.18.6(2/2)
- Ordered_Sets
 - child of* Ada.Containers A.18.9(2/2)
- Pointers
 - child of* Interfaces.C B.3.2(4)
- Real_Arrays
 - child of* Ada.Numerics G.3.1(31/2)
- Real_Time
 - child of* Ada D.8(3)
- Round_Robin
 - child of* Ada.Dispatching D.2.5(4/2)
- RPC
 - child of* System E.5(3)
- Sequential_IO
 - child of* Ada A.8.1(2)
- Single_Precision_Complex_Types
 - in* Interfaces.Fortran B.5(8)
- Standard A.1(4)
- Storage_Elements
 - child of* System 13.7.1(2/2)
- Storage_IO
 - child of* Ada A.9(3)
- Storage_Pools
 - child of* System 13.11(5)
- Stream_IO
 - child of* Ada.Streams A.12.1(3)
- Streams
 - child of* Ada 13.13.1(2)
- Strings
 - child of* Ada A.4.1(3)
 - child of* Interfaces.C B.3.1(3)
- Synchronous_Task_Control
 - child of* Ada D.10(3/2)
- System 13.7(3/2)

Tags	
<i>child of Ada</i> 3.9(6/2)	
Task_Attributes	
<i>child of Ada</i> C.7.2(2)	
Task_Identification	
<i>child of Ada</i> C.7.1(2/2)	
Task_Termination	
<i>child of Ada</i> C.7.3(2/2)	
Text_Streams	
<i>child of Ada.Text_IO</i> A.12.2(3)	
<i>child of Ada.Wide_Text_IO</i> A.12.3(3)	
<i>child of Ada.Wide_Wide_Text_IO</i> A.12.4(3/2)	
Text_IO	
<i>child of Ada</i> A.10.1(2)	
Time_Zones	
<i>child of Ada.Calendar</i> 9.6.1(2/2)	
Timers	
<i>child of Ada.Execution_Time</i> D.14.1(3/2)	
Timing_Events	
<i>child of Ada.Real_Time</i> D.15(3/2)	
Unbounded	
<i>child of Ada.Strings</i> A.4.5(3)	
Unbounded_IO	
<i>child of Ada.Text_IO</i> A.10.12(3/2)	
<i>child of Ada.Wide_Text_IO</i> A.11(5/2)	
<i>child of Ada.Wide_Wide_Text_IO</i> A.11(5/2)	
Vectors	
<i>child of Ada.Containers</i> A.18.2(6/2)	
Wide_Bounded	
<i>child of Ada.Strings</i> A.4.7(1/2)	
Wide_Constants	
<i>child of Ada.Strings.Wide_Maps</i> A.4.7(1/2), A.4.8(28/2)	
Wide_Fixed	
<i>child of Ada.Strings</i> A.4.7(1/2)	
Wide_Hash	
<i>child of Ada.Strings</i> A.4.7(1/2)	
Wide_Maps	
<i>child of Ada.Strings</i> A.4.7(3)	
Wide_Text_IO	
<i>child of Ada</i> A.11(2/2)	
Wide_Unbounded	
<i>child of Ada.Strings</i> A.4.7(1/2)	
Wide_Characters	
<i>child of Ada</i> A.3.1(4/2)	
Wide_Wide_Constants	
<i>child of Ada.Strings.Wide_Wide_Maps</i> A.4.8(1/2)	
Wide_Wide_Hash	
<i>child of Ada.Strings</i> A.4.8(1/2)	
Wide_Wide_Text_IO	
<i>child of Ada</i> A.11(3/2)	
Wide_Wide_Bounded	
<i>child of Ada.Strings</i> A.4.8(1/2)	
Wide_Wide_Characters	
<i>child of Ada</i> A.3.1(6/2)	
Wide_Wide_Fixed	
<i>child of Ada.Strings</i> A.4.8(1/2)	
Wide_Wide_Maps	
<i>child of Ada.Strings</i> A.4.8(3/2)	
Wide_Wide_Unbounded	
<i>child of Ada.Strings</i> A.4.8(1/2)	

Q.2 Language-Defined Types and Subtypes

This clause lists all language-defined types and subtypes.

1/2

Address	
<i>in System</i> 13.7(12)	
Alignment	
<i>in Ada.Strings</i> A.4.1(6)	
Alphanumeric	
<i>in Interfaces.COBOL</i> B.4(16)	
Any_Priority subtype of Integer	
<i>in System</i> 13.7(16)	
Attribute_Handle	
<i>in Ada.Task_Attributes</i> C.7.2(3)	
Binary	
<i>in Interfaces.COBOL</i> B.4(10)	
Binary_Format	
<i>in Interfaces.COBOL</i> B.4(24)	
Bit_Order	
<i>in System</i> 13.7(15/2)	
Boolean	
<i>in Standard</i> A.1(5)	
Bounded_String	
<i>in Ada.Strings.Bounded</i> A.4.4(6)	
Buffer_Type subtype of Storage_Array	
<i>in Ada.Storage_IO</i> A.9(4)	
Byte	
<i>in Interfaces.COBOL</i> B.4(29)	
Byte_Array	
<i>in Interfaces.COBOL</i> B.4(29)	
C_float	
<i>in Interfaces.C</i> B.3(15)	
Cause_Of_Termination	
<i>in Ada.Task_Termination</i> C.7.3(3/2)	
char	
<i>in Interfaces.C</i> B.3(19)	
char16_array	
<i>in Interfaces.C</i> B.3(39.5/2)	
char16_t	
<i>in Interfaces.C</i> B.3(39.2/2)	
char32_array	
<i>in Interfaces.C</i> B.3(39.14/2)	
char32_t	
<i>in Interfaces.C</i> B.3(39.11/2)	

- char_array
 - in Interfaces.C* B.3(23)
- char_array_access
 - in Interfaces.C.Strings* B.3.1(4)
- Character
 - in Standard* A.1(35/2)
- Character_Mapping
 - in Ada.Strings.Maps* A.4.2(20/2)
- Character_Mapping_Function
 - in Ada.Strings.Maps* A.4.2(25)
- Character_Range
 - in Ada.Strings.Maps* A.4.2(6)
- Character_Ranges
 - in Ada.Strings.Maps* A.4.2(7)
- Character_Sequence *subtype of String*
 - in Ada.Strings.Maps* A.4.2(16)
- Character_Set
 - in Ada.Strings.Maps* A.4.2(4/2)
 - in Interfaces.Fortran* B.5(11)
- chars_ptr
 - in Interfaces.C.Strings* B.3.1(5/2)
- chars_ptr_array
 - in Interfaces.C.Strings* B.3.1(6/2)
- COBOL_Character
 - in Interfaces.COBOL* B.4(13)
- Complex
 - in Ada.Numerics.Generic_Complex_Types* G.1.1(3)
 - in Interfaces.Fortran* B.5(9)
- Complex_Matrix
 - in Ada.Numerics.Generic_Complex_Arrays* G.3.2(4/2)
- Complex_Vector
 - in Ada.Numerics.Generic_Complex_Arrays* G.3.2(4/2)
- Controlled
 - in Ada.Finalization* 7.6(5/2)
- Count
 - in Ada.Direct_IO* A.8.4(4)
 - in Ada.Streams.Stream_IO* A.12.1(7)
 - in Ada.Text_IO* A.10.1(5)
- CPU_Time
 - in Ada.Execution_Time* D.14(4/2)
- Cursor
 - in Ada.Containers.Doubly_Linked_Lists* A.18.3(7/2)
 - in Ada.Containers.Hashed_Maps* A.18.5(4/2)
 - in Ada.Containers.Hashed_Sets* A.18.8(4/2)
 - in Ada.Containers.Ordered_Maps* A.18.6(5/2)
 - in Ada.Containers.Ordered_Sets* A.18.9(5/2)
 - in Ada.Containers.Vectors* A.18.2(9/2)
- Day_Count
 - in Ada.Calendar.Arithmetic* 9.6.1(10/2)
- Day_Duration *subtype of Duration*
 - in Ada.Calendar* 9.6(11/2)
- Day_Name
 - in Ada.Calendar.Formatting* 9.6.1(17/2)
- Day_Number *subtype of Integer*
 - in Ada.Calendar* 9.6(11/2)
- Deadline *subtype of Time*
 - in Ada.Dispatching.EDF* D.2.6(9/2)
- Decimal_Element
 - in Interfaces.COBOL* B.4(12)
- Direction
 - in Ada.Strings* A.4.1(6)
- Directory_Entry_Type
 - in Ada.Directories* A.16(29/2)
- Display_Format
 - in Interfaces.COBOL* B.4(22)
- double
 - in Interfaces.C* B.3(16)
- Double_Precision
 - in Interfaces.Fortran* B.5(6)
- Duration
 - in Standard* A.1(43)
- Exception_Id
 - in Ada.Exceptions* 11.4.1(2/2)
- Exception_Occurrence
 - in Ada.Exceptions* 11.4.1(3/2)
- Exception_Occurrence_Access
 - in Ada.Exceptions* 11.4.1(3/2)
- Exit_Status
 - in Ada.Command_Line* A.15(7)
- Extended_Index *subtype of Index_Type'Base*
 - in Ada.Containers.Vectors* A.18.2(7/2)
- Field *subtype of Integer*
 - in Ada.Text_IO* A.10.1(6)
- File_Access
 - in Ada.Text_IO* A.10.1(18)
- File_Kind
 - in Ada.Directories* A.16(22/2)
- File_Mode
 - in Ada.Direct_IO* A.8.4(4)
 - in Ada.Sequential_IO* A.8.1(4)
 - in Ada.Streams.Stream_IO* A.12.1(6)
 - in Ada.Text_IO* A.10.1(4)
- File_Size
 - in Ada.Directories* A.16(23/2)
- File_Type
 - in Ada.Direct_IO* A.8.4(3)
 - in Ada.Sequential_IO* A.8.1(3)
 - in Ada.Streams.Stream_IO* A.12.1(5)
 - in Ada.Text_IO* A.10.1(3)
- Filter_Type
 - in Ada.Directories* A.16(30/2)
- Float
 - in Standard* A.1(21)
- Floating
 - in Interfaces.COBOL* B.4(9)
- Fortran_Character
 - in Interfaces.Fortran* B.5(12)
- Fortran_Integer
 - in Interfaces.Fortran* B.5(5)
- Generator
 - in Ada.Numerics.Discrete_Random* A.5.2(19)
 - in Ada.Numerics.Float_Random* A.5.2(7)
- Group_Budget
 - in Ada.Execution_Time.Group_Budgets* D.14.2(4/2)
- Group_Budget_Handler
 - in Ada.Execution_Time.Group_Budgets* D.14.2(5/2)
- Hash_Type
 - in Ada.Containers* A.18.1(4/2)

Hour_Number *subtype of* Natural
in Ada.Calendar.Formatting 9.6.1(20/2)

Imaginary
in Ada.Numerics.Generic_Complex_Types G.1.1(4/2)

Imaginary *subtype of* Imaginary
in Interfaces.Fortran B.5(10)

int
in Interfaces.C B.3(7)

Integer
in Standard A.1(12)

Integer_Address
in System.Storage_Elements 13.7.1(10)

Interrupt_ID
in Ada.Interrupts C.3.2(2)

Interrupt_Priority *subtype of* Any_Priority
in System 13.7(16)

ISO_646 *subtype of* Character
in Ada.Characters.Handling A.3.2(9)

Leap_Seconds_Count *subtype of* Integer
in Ada.Calendar.Arithmetic 9.6.1(11/2)

Length_Range *subtype of* Natural
in Ada.Strings.Bounded A.4.4(8)

Limited_Controlled
in Ada.Finalization 7.6(7/2)

List
in Ada.Containers.Doubly_Linked_Lists A.18.3(6/2)

Logical
in Interfaces.Fortran B.5(7)

long
in Interfaces.C B.3(7)

Long_Binary
in Interfaces.COBOL B.4(10)

long_double
in Interfaces.C B.3(17)

Long_Floating
in Interfaces.COBOL B.4(9)

Map
in Ada.Containers.Hashed_Maps A.18.5(3/2)
in Ada.Containers.Ordered_Maps A.18.6(4/2)

Membership
in Ada.Strings A.4.1(6)

Minute_Number *subtype of* Natural
in Ada.Calendar.Formatting 9.6.1(20/2)

Month_Number *subtype of* Integer
in Ada.Calendar 9.6(11/2)

Name
in System 13.7(4)

Natural *subtype of* Integer
in Standard A.1(13)

Number_Base *subtype of* Integer
in Ada.Text_IO A.10.1(6)

Numeric
in Interfaces.COBOL B.4(20)

Packed_Decimal
in Interfaces.COBOL B.4(12)

Packed_Format
in Interfaces.COBOL B.4(26)

Parameterless_Handler
in Ada.Interrupts C.3.2(2)

Params_Stream_Type
in System.RPC E.5(6)

Partition_Id
in System.RPC E.5(4)

Picture
in Ada.Text_IO.Editing F.3.3(4)

plain_char
in Interfaces.C B.3(11)

Pointer
in Interfaces.C.Pointers B.3.2(5)

Positive *subtype of* Integer
in Standard A.1(13)

Positive_Count *subtype of* Count
in Ada.Direct_IO A.8.4(4)
in Ada.Streams.Stream_IO A.12.1(7)
in Ada.Text_IO A.10.1(5)

Priority *subtype of* Any_Priority
in System 13.7(16)

ptrdiff_t
in Interfaces.C B.3(12)

Real
in Interfaces.Fortran B.5(6)

Real_Matrix
in Ada.Numerics.Generic_Real_Arrays G.3.1(4/2)

Real_Vector
in Ada.Numerics.Generic_Real_Arrays G.3.1(4/2)

Root_Storage_Pool
in System.Storage_Pools 13.11(6/2)

Root_Stream_Type
in Ada.Streams 13.13.1(3/2)

RPC_Receiver
in System.RPC E.5(11)

Search_Type
in Ada.Directories A.16(31/2)

Second_Duration *subtype of* Day_Duration
in Ada.Calendar.Formatting 9.6.1(20/2)

Second_Number *subtype of* Natural
in Ada.Calendar.Formatting 9.6.1(20/2)

Seconds_Count
in Ada.Real_Time D.8(15)

Set
in Ada.Containers.Hashed_Sets A.18.8(3/2)
in Ada.Containers.Ordered_Sets A.18.9(4/2)

short
in Interfaces.C B.3(7)

signed_char
in Interfaces.C B.3(8)

size_t
in Interfaces.C B.3(13)

State
in Ada.Numerics.Discrete_Random A.5.2(23)
in Ada.Numerics.Float_Random A.5.2(11)

Storage_Array
in System.Storage_Elements 13.7.1(5)

Storage_Count *subtype of* Storage_Offset
in System.Storage_Elements 13.7.1(4)

Storage_Element
in System.Storage_Elements 13.7.1(5)

Storage_Offset
in System.Storage_Elements 13.7.1(3)

Stream_Access
in Ada.Streams.Stream_IO A.12.1(4)
in Ada.Text_IO.Text_Streams A.12.2(3)
in Ada.Wide_Text_IO.Text_Streams A.12.3(3)
in Ada.Wide_Wide_Text_IO.Text_Streams A.12.4(3/2)

Stream_Element
in Ada.Streams 13.13.1(4/1)

Stream_Element_Array
in Ada.Streams 13.13.1(4/1)

Stream_Element_Count *subtype of Stream_Element_Offset*
in Ada.Streams 13.13.1(4/1)

Stream_Element_Offset
in Ada.Streams 13.13.1(4/1)

String
in Standard A.1(37)

String_Access
in Ada.Strings.Unbounded A.4.5(7)

Suspension_Object
in Ada.Synchronous_Task_Control D.10(4)

Tag
in Ada.Tags 3.9(6/2)

Tag_Array
in Ada.Tags 3.9(7.3/2)

Task_Array
in Ada.Execution_Time.Group_Budgets D.14.2(6/2)

Task_Id
in Ada.Task_Identification C.7.1(2/2)

Termination_Handler
in Ada.Task_Termination C.7.3(4/2)

Time
in Ada.Calendar 9.6(10)
in Ada.Real_Time D.8(4)

Time_Offset
in Ada.Calendar.Time_Zones 9.6.1(4/2)

Time_Span
in Ada.Real_Time D.8(5)

Timer
in Ada.Execution_Time.Timers D.14.1(4/2)

Timer_Handler
in Ada.Execution_Time.Timers D.14.1(5/2)

Timing_Event
in Ada.Real_Time.Timing_Events D.15(4/2)

Timing_Event_Handler
in Ada.Real_Time.Timing_Events D.15(4/2)

Trim_End
in Ada.Strings A.4.1(6)

Truncation
in Ada.Strings A.4.1(6)

Type_Set
in Ada.Text_IO A.10.1(7)

Unbounded_String
in Ada.Strings.Unbounded A.4.5(4/2)

Uniformly_Distributed *subtype of Float*
in Ada.Numerics.Float_Random A.5.2(8)

unsigned
in Interfaces.C B.3(9)

unsigned_char
in Interfaces.C B.3(10)

unsigned_long
in Interfaces.C B.3(9)

unsigned_short
in Interfaces.C B.3(9)

Vector
in Ada.Containers.Vectors A.18.2(8/2)

wchar_array
in Interfaces.C B.3(33)

wchar_t
in Interfaces.C B.3(30/1)

Wide_Character
in Standard A.1(36.1/2)

Wide_Character_Mapping
in Ada.Strings.Wide_Maps A.4.7(20/2)

Wide_Character_Mapping_Function
in Ada.Strings.Wide_Maps A.4.7(26)

Wide_Character_Range
in Ada.Strings.Wide_Maps A.4.7(6)

Wide_Character_Ranges
in Ada.Strings.Wide_Maps A.4.7(7)

Wide_Character_Sequence *subtype of Wide_String*
in Ada.Strings.Wide_Maps A.4.7(16)

Wide_Character_Set
in Ada.Strings.Wide_Maps A.4.7(4/2)

Wide_String
in Standard A.1(41)

Wide_Wide_Character
in Standard A.1(36.2/2)

Wide_Wide_Character_Mapping
in Ada.Strings.Wide_Wide_Maps A.4.8(20/2)

Wide_Wide_Character_Mapping_Function
in Ada.Strings.Wide_Wide_Maps A.4.8(26/2)

Wide_Wide_Character_Range
in Ada.Strings.Wide_Wide_Maps A.4.8(6/2)

Wide_Wide_Character_Ranges
in Ada.Strings.Wide_Wide_Maps A.4.8(7/2)

Wide_Wide_Character_Sequence *subtype of Wide_Wide_String*
in Ada.Strings.Wide_Wide_Maps A.4.8(16/2)

Wide_Wide_Character_Set
in Ada.Strings.Wide_Wide_Maps A.4.8(4/2)

Wide_Wide_String
in Standard A.1(42.1/2)

Year_Number *subtype of Integer*
in Ada.Calendar 9.6(11/2)

Q.3 Language-Defined Subprograms

This clause lists all language-defined subprograms.

1/2

- Abort_Task *in Ada.Task_Identification* C.7.1(3/1)
- Actual_Quantum
 - in Ada.Dispatching.Round_Robin* D.2.5(4/2)
- Add
 - in Ada.Execution_Time.Group_Budgets* D.14.2(9/2)
- Add_Task
 - in Ada.Execution_Time.Group_Budgets* D.14.2(8/2)
- Adjust *in Ada.Finalization* 7.6(6/2)
- Allocate *in System.Storage_Pools* 13.11(7)
- Append
 - in Ada.Containers.Doubly_Linked_Lists* A.18.3(23/2)
 - in Ada.Containers.Vectors* A.18.2(46/2), A.18.2(47/2)
 - in Ada.Strings.Bounded* A.4.4(13), A.4.4(14), A.4.4(15), A.4.4(16), A.4.4(17), A.4.4(18), A.4.4(19), A.4.4(20)
 - in Ada.Strings.Unbounded* A.4.5(12), A.4.5(13), A.4.5(14)
- Arccos
 - in Ada.Numerics.Generic_Complex_Elementary_Functions* G.1.2(5)
 - in Ada.Numerics.Generic_Elementary_Functions* A.5.1(6)
- Arccosh
 - in Ada.Numerics.Generic_Complex_Elementary_Functions* G.1.2(7)
 - in Ada.Numerics.Generic_Elementary_Functions* A.5.1(7)
- Arccot
 - in Ada.Numerics.Generic_Complex_Elementary_Functions* G.1.2(5)
 - in Ada.Numerics.Generic_Elementary_Functions* A.5.1(6)
- Arccoth
 - in Ada.Numerics.Generic_Complex_Elementary_Functions* G.1.2(7)
 - in Ada.Numerics.Generic_Elementary_Functions* A.5.1(7)
- Arcsin
 - in Ada.Numerics.Generic_Complex_Elementary_Functions* G.1.2(5)
 - in Ada.Numerics.Generic_Elementary_Functions* A.5.1(6)
- Arcsinh
 - in Ada.Numerics.Generic_Complex_Elementary_Functions* G.1.2(7)
 - in Ada.Numerics.Generic_Elementary_Functions* A.5.1(7)
- Arctan
 - in Ada.Numerics.Generic_Complex_Elementary_Functions* G.1.2(5)
 - in Ada.Numerics.Generic_Elementary_Functions* A.5.1(6)
- Arctanh
 - in Ada.Numerics.Generic_Complex_Elementary_Functions* G.1.2(7)
 - in Ada.Numerics.Generic_Elementary_Functions* A.5.1(7)
- Argument
 - in Ada.Command_Line* A.15(5)
 - in Ada.Numerics.Generic_Complex_Arrays* G.3.2(10/2), G.3.2(31/2)
 - in Ada.Numerics.Generic_Complex_Types* G.1.1(10)
- Argument_Count *in Ada.Command_Line* A.15(4)
- Attach_Handler *in Ada.Interrupts* C.3.2(7)
- Base_Name *in Ada.Directories* A.16(19/2)
- Blank_When_Zero
 - in Ada.Text_IO.Editing* F.3.3(7)
- Bounded_Slice *in Ada.Strings.Bounded* A.4.4(28.1/2), A.4.4(28.2/2)
- Budget_Has_Expired
 - in Ada.Execution_Time.Group_Budgets* D.14.2(9/2)
- Budget_Remaining
 - in Ada.Execution_Time.Group_Budgets* D.14.2(9/2)
- Cancel_Handler
 - in Ada.Execution_Time.Group_Budgets* D.14.2(10/2)
 - in Ada.Execution_Time.Timers* D.14.1(7/2)
 - in Ada.Real_Time.Timing_Events* D.15(5/2)
- Capacity
 - in Ada.Containers.Hashed_Maps* A.18.5(8/2)
 - in Ada.Containers.Hashed_Sets* A.18.8(10/2)
 - in Ada.Containers.Vectors* A.18.2(19/2)
- Ceiling
 - in Ada.Containers.Ordered_Maps* A.18.6(41/2)
 - in Ada.Containers.Ordered_Sets* A.18.9(51/2), A.18.9(71/2)
- Clear
 - in Ada.Containers.Doubly_Linked_Lists* A.18.3(13/2)
 - in Ada.Containers.Hashed_Maps* A.18.5(12/2)
 - in Ada.Containers.Hashed_Sets* A.18.8(14/2)
 - in Ada.Containers.Ordered_Maps* A.18.6(11/2)
 - in Ada.Containers.Ordered_Sets* A.18.9(13/2)
 - in Ada.Containers.Vectors* A.18.2(24/2)
 - in Ada.Environment_Variables* A.17(7/2)
- Clock
 - in Ada.Calendar* 9.6(12)
 - in Ada.Execution_Time* D.14(5/2)
 - in Ada.Real_Time* D.8(6)
- Close
 - in Ada.Direct_IO* A.8.4(8)
 - in Ada.Sequential_IO* A.8.1(8)
 - in Ada.Streams.Stream_IO* A.12.1(10)
 - in Ada.Text_IO* A.10.1(11)
- Col *in Ada.Text_IO* A.10.1(37)
- Command_Name *in Ada.Command_Line* A.15(6)
- Compose *in Ada.Directories* A.16(20/2)
- Compose_From_Cartesian
 - in Ada.Numerics.Generic_Complex_Arrays* G.3.2(9/2), G.3.2(29/2)
 - in Ada.Numerics.Generic_Complex_Types* G.1.1(8)
- Compose_From_Polar
 - in Ada.Numerics.Generic_Complex_Arrays* G.3.2(11/2), G.3.2(32/2)
 - in Ada.Numerics.Generic_Complex_Types* G.1.1(11)

Conjugate
in Ada.Numerics.Generic_Complex_Arrays G.3.2(13/2),
 G.3.2(34/2)
in Ada.Numerics.Generic_Complex_Types G.1.1(12),
 G.1.1(15)

Containing_Directory
in Ada.Directories A.16(17/2)

Contains
in Ada.Containers.Doubly_Linked_Lists A.18.3(43/2)
in Ada.Containers.Hashed_Maps A.18.5(32/2)
in Ada.Containers.Hashed_Sets A.18.8(44/2), A.18.8(57/2)
in Ada.Containers.Ordered_Maps A.18.6(42/2)
in Ada.Containers.Ordered_Sets A.18.9(52/2), A.18.9(72/2)
in Ada.Containers.Vectors A.18.2(71/2)

Continue
in Ada.Aynchronous_Task_Control D.11(3/2)

Copy_Array *in Interfaces.C.Pointers* B.3.2(15)

Copy_File *in Ada.Directories* A.16(13/2)

Copy_Terminated_Array
in Interfaces.C.Pointers B.3.2(14)

Cos
in Ada.Numerics.Generic_Complex_Elementary_Functions
 G.1.2(4)
in Ada.Numerics.Generic_Elementary_Functions A.5.1(5)

Cosh
in Ada.Numerics.Generic_Complex_Elementary_Functions
 G.1.2(6)
in Ada.Numerics.Generic_Elementary_Functions A.5.1(7)

Cot
in Ada.Numerics.Generic_Complex_Elementary_Functions
 G.1.2(4)
in Ada.Numerics.Generic_Elementary_Functions A.5.1(5)

Coth
in Ada.Numerics.Generic_Complex_Elementary_Functions
 G.1.2(6)
in Ada.Numerics.Generic_Elementary_Functions A.5.1(7)

Count
in Ada.Strings.Bounded A.4.4(48), A.4.4(49), A.4.4(50)
in Ada.Strings.Fixed A.4.3(13), A.4.3(14), A.4.3(15)
in Ada.Strings.Unbounded A.4.5(43), A.4.5(44), A.4.5(45)

Create
in Ada.Direct_IO A.8.4(6)
in Ada.Sequential_IO A.8.1(6)
in Ada.Streams.Stream_IO A.12.1(8)
in Ada.Text_IO A.10.1(9)

Create_Directory *in Ada.Directories* A.16(7/2)

Create_Path *in Ada.Directories* A.16(9/2)

Current_Directory *in Ada.Directories* A.16(5/2)

Current_Error *in Ada.Text_IO* A.10.1(17), A.10.1(20)

Current_Handler
in Ada.Execution_Time.Group_Budgets D.14.2(10/2)
in Ada.Execution_Time.Timers D.14.1(7/2)
in Ada.Interrupts C.3.2(6)
in Ada.Real_Time.Timing_Events D.15(5/2)

Current_Input *in Ada.Text_IO* A.10.1(17), A.10.1(20)

Current_Output *in Ada.Text_IO* A.10.1(17), A.10.1(20)

Current_State
in Ada.Synchronous_Task_Control D.10(4)

Current_Task
in Ada.Task_Identification C.7.1(3/1)

Current_Task_Fallback_Handler
in Ada.Task_Termination C.7.3(5/2)

Day
in Ada.Calendar 9.6(13)
in Ada.Calendar.Formatting 9.6.1(23/2)

Day_of_Week
in Ada.Calendar.Formatting 9.6.1(18/2)

Deallocate *in System.Storage_Pools* 13.11(8)

Decrement *in Interfaces.C.Pointers* B.3.2(11)

Delay_Until_And_Set_Deadline
in Ada.Dispatching.EDF D.2.6(9/2)

Delete
in Ada.Containers.Doubly_Linked_Lists A.18.3(24/2)
in Ada.Containers.Hashed_Maps A.18.5(25/2), A.18.5(26/2)
in Ada.Containers.Hashed_Sets A.18.8(24/2), A.18.8(25/2),
 A.18.8(55/2)
in Ada.Containers.Ordered_Maps A.18.6(24/2),
 A.18.6(25/2)
in Ada.Containers.Ordered_Sets A.18.9(23/2), A.18.9(24/2),
 A.18.9(68/2)
in Ada.Containers.Vectors A.18.2(50/2), A.18.2(51/2)
in Ada.Direct_IO A.8.4(8)
in Ada.Sequential_IO A.8.1(8)
in Ada.Streams.Stream_IO A.12.1(10)
in Ada.Strings.Bounded A.4.4(64), A.4.4(65)
in Ada.Strings.Fixed A.4.3(29), A.4.3(30)
in Ada.Strings.Unbounded A.4.5(59), A.4.5(60)
in Ada.Text_IO A.10.1(11)

Delete_Directory *in Ada.Directories* A.16(8/2)

Delete_File *in Ada.Directories* A.16(11/2)

Delete_First
in Ada.Containers.Doubly_Linked_Lists A.18.3(25/2)
in Ada.Containers.Ordered_Maps A.18.6(26/2)
in Ada.Containers.Ordered_Sets A.18.9(25/2)
in Ada.Containers.Vectors A.18.2(52/2)

Delete_Last
in Ada.Containers.Doubly_Linked_Lists A.18.3(26/2)
in Ada.Containers.Ordered_Maps A.18.6(27/2)
in Ada.Containers.Ordered_Sets A.18.9(26/2)
in Ada.Containers.Vectors A.18.2(53/2)

Delete_Tree *in Ada.Directories* A.16(10/2)

Dereference_Error
in Interfaces.C.Strings B.3.1(12)

Descendant_Tag *in Ada.Tags* 3.9(7.1/2)

Detach_Handler *in Ada.Interrupts* C.3.2(9)

Determinant
in Ada.Numerics.Generic_Complex_Arrays G.3.2(46/2)
in Ada.Numerics.Generic_Real_Arrays G.3.1(24/2)

Difference
in Ada.Calendar.Arithmetic 9.6.1(12/2)
in Ada.Containers.Hashed_Sets A.18.8(32/2), A.18.8(33/2)
in Ada.Containers.Ordered_Sets A.18.9(33/2), A.18.9(34/2)

Divide *in Ada.Decimal* F.2(6)

Do_APIC *in System.RPC* E.5(10)

Do_RPC *in System.RPC* E.5(9)

Eigenystem
in Ada.Numerics.Generic_Complex_Arrays G.3.2(49/2)
in Ada.Numerics.Generic_Real_Arrays G.3.1(27/2)

Eigenvalues
in Ada.Numerics.Generic_Complex_Arrays G.3.2(48/2)
in Ada.Numerics.Generic_Real_Arrays G.3.1(26/2)

Element
in Ada.Containers.Doubly_Linked_Lists A.18.3(14/2)
in Ada.Containers.Hashed_Maps A.18.5(14/2), A.18.5(31/2)
in Ada.Containers.Hashed_Sets A.18.8(15/2), A.18.8(52/2)
in Ada.Containers.Ordered_Maps A.18.6(13/2),
A.18.6(39/2)
in Ada.Containers.Ordered_Sets A.18.9(14/2), A.18.9(65/2)
in Ada.Containers.Vectors A.18.2(27/2), A.18.2(28/2)
in Ada.Strings.Bounded A.4.4(26)
in Ada.Strings.Unbounded A.4.5(20)

End_Of_File
in Ada.Direct_IO A.8.4(16)
in Ada.Sequential_IO A.8.1(13)
in Ada.Streams.Stream_IO A.12.1(12)
in Ada.Text_IO A.10.1(34)

End_Of_Line *in Ada.Text_IO* A.10.1(30)

End_Of_Page *in Ada.Text_IO* A.10.1(33)

End_Search *in Ada.Directories* A.16(33/2)

Equivalent_Elements
in Ada.Containers.Hashed_Sets A.18.8(46/2), A.18.8(47/2),
A.18.8(48/2)
in Ada.Containers.Ordered_Sets A.18.9(3/2)

Equivalent_Keys
in Ada.Containers.Hashed_Maps A.18.5(34/2),
A.18.5(35/2), A.18.5(36/2)
in Ada.Containers.Ordered_Maps A.18.6(3/2)
in Ada.Containers.Ordered_Sets A.18.9(63/2)

Equivalent_Sets
in Ada.Containers.Hashed_Sets A.18.8(8/2)
in Ada.Containers.Ordered_Sets A.18.9(9/2)

Establish_RPC_Receiver *in System.RPC* E.5(12)

Exception_Identity *in Ada.Exceptions* 11.4.1(5/2)

Exception_Information
in Ada.Exceptions 11.4.1(5/2)

Exception_Message *in Ada.Exceptions* 11.4.1(4/2)

Exception_Name *in Ada.Exceptions* 11.4.1(2/2), 11.4.1(5/2)

Exchange_Handler *in Ada.Interrupts* C.3.2(8)

Exclude
in Ada.Containers.Hashed_Maps A.18.5(24/2)
in Ada.Containers.Hashed_Sets A.18.8(23/2), A.18.8(54/2)
in Ada.Containers.Ordered_Maps A.18.6(23/2)
in Ada.Containers.Ordered_Sets A.18.9(22/2), A.18.9(67/2)

Exists
in Ada.Directories A.16(24/2)
in Ada.Environment_Variables A.17(5/2)

Exp
in Ada.Numerics.Generic_Complex_Elementary_Functions
G.1.2(3)
in Ada.Numerics.Generic_Elementary_Functions A.5.1(4)

Expanded_Name *in Ada.Tags* 3.9(7/2)

Extension *in Ada.Directories* A.16(18/2)

External_Tag *in Ada.Tags* 3.9(7/2)

Finalize *in Ada.Finalization* 7.6(6/2), 7.6(8/2)

Find
in Ada.Containers.Doubly_Linked_Lists A.18.3(41/2)
in Ada.Containers.Hashed_Maps A.18.5(30/2)
in Ada.Containers.Hashed_Sets A.18.8(43/2), A.18.8(56/2)

in Ada.Containers.Ordered_Maps A.18.6(38/2)
in Ada.Containers.Ordered_Sets A.18.9(49/2), A.18.9(69/2)
in Ada.Containers.Vectors A.18.2(68/2)

Find_Index *in Ada.Containers.Vectors* A.18.2(67/2)

Find_Token
in Ada.Strings.Bounded A.4.4(51)
in Ada.Strings.Fixed A.4.3(16)
in Ada.Strings.Unbounded A.4.5(46)

First
in Ada.Containers.Doubly_Linked_Lists A.18.3(33/2)
in Ada.Containers.Hashed_Maps A.18.5(27/2)
in Ada.Containers.Hashed_Sets A.18.8(40/2)
in Ada.Containers.Ordered_Maps A.18.6(28/2)
in Ada.Containers.Ordered_Sets A.18.9(41/2)
in Ada.Containers.Vectors A.18.2(58/2)

First_Element
in Ada.Containers.Doubly_Linked_Lists A.18.3(34/2)
in Ada.Containers.Ordered_Maps A.18.6(29/2)
in Ada.Containers.Ordered_Sets A.18.9(42/2)
in Ada.Containers.Vectors A.18.2(59/2)

First_Index *in Ada.Containers.Vectors* A.18.2(57/2)

First_Key
in Ada.Containers.Ordered_Maps A.18.6(30/2)

Floor
in Ada.Containers.Ordered_Maps A.18.6(40/2)
in Ada.Containers.Ordered_Sets A.18.9(50/2), A.18.9(70/2)

Flush
in Ada.Streams.Stream_IO A.12.1(25/1)
in Ada.Text_IO A.10.1(21/1)

Form
in Ada.Direct_IO A.8.4(9)
in Ada.Sequential_IO A.8.1(9)
in Ada.Streams.Stream_IO A.12.1(11)
in Ada.Text_IO A.10.1(12)

Free
in Ada.Strings.Unbounded A.4.5(7)
in Interfaces.C.Strings B.3.1(11)

Full_Name *in Ada.Directories* A.16(15/2), A.16(39/2)

Generic_Array_Sort
child of Ada.Containers A.18.16(3/2)

Generic_Constrained_Array_Sort
child of Ada.Containers A.18.16(7/2)

Get
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in Ada.Numerics.Generic_Complex_Elementary_Functions
 G.1.2(6)
in Ada.Numerics.Generic_Elementary_Functions A.5.1(7)
 Size
in Ada.Direct_IO A.8.4(15)
in Ada.Directories A.16(26/2), A.16(41/2)
in Ada.Streams.Stream_IO A.12.1(23)
 Skip_Line *in* Ada.Text_IO A.10.1(29)
 Skip_Page *in* Ada.Text_IO A.10.1(32)
 Slice
in Ada.Strings.Bounded A.4.4(28)
in Ada.Strings.Unbounded A.4.5(22)
 Solve
in Ada.Numerics.Generic_Complex_Arrays G.3.2(46/2)
in Ada.Numerics.Generic_Real_Arrays G.3.1(24/2)
 Sort
in Ada.Containers.Doubly_Linked_Lists A.18.3(49/2)
in Ada.Containers.Vectors A.18.2(77/2)
 Specific_Handler
in Ada.Task_Termination C.7.3(6/2)
 Splice
in Ada.Containers.Doubly_Linked_Lists A.18.3(30/2),
 A.18.3(31/2), A.18.3(32/2)
 Split
in Ada.Calendar 9.6(14)
in Ada.Calendar.Formatting 9.6.1(29/2), 9.6.1(32/2),
 9.6.1(33/2), 9.6.1(34/2)
in Ada.Execution_Time D.14(8/2)
in Ada.Real_Time D.8(16)
 Sqrt
in Ada.Numerics.Generic_Complex_Elementary_Functions
 G.1.2(3)
in Ada.Numerics.Generic_Elementary_Functions A.5.1(4)
 Standard_Error *in* Ada.Text_IO A.10.1(16), A.10.1(19)
 Standard_Input *in* Ada.Text_IO A.10.1(16), A.10.1(19)
 Standard_Output *in* Ada.Text_IO A.10.1(16), A.10.1(19)
 Start_Search *in* Ada.Directories A.16(32/2)
 Storage_Size *in* System.Storage_Pools 13.11(9)
 Stream
in Ada.Streams.Stream_IO A.12.1(13)
in Ada.Text_IO.Text_Streams A.12.2(4)
in Ada.Wide_Text_IO.Text_Streams A.12.3(4)
in Ada.Wide_Wide_Text_IO.Text_Streams A.12.4(4/2)
 Strlen *in* Interfaces.C.Strings B.3.1(17)
 Sub_Second *in* Ada.Calendar.Formatting 9.6.1(27/2)
 Suspend_Until_True
in Ada.Synchronous_Task_Control D.10(4)
 Swap
in Ada.Containers.Doubly_Linked_Lists A.18.3(28/2)
in Ada.Containers.Vectors A.18.2(55/2), A.18.2(56/2)
 Swap_Links
in Ada.Containers.Doubly_Linked_Lists A.18.3(29/2)
 Symmetric_Difference
in Ada.Containers.Hashed_Sets A.18.8(35/2), A.18.8(36/2)
in Ada.Containers.Ordered_Sets A.18.9(36/2), A.18.9(37/2)
 Tail
in Ada.Strings.Bounded A.4.4(72), A.4.4(73)
in Ada.Strings.Fixed A.4.3(37), A.4.3(38)
in Ada.Strings.Unbounded A.4.5(67), A.4.5(68)
 Tan
in Ada.Numerics.Generic_Complex_Elementary_Functions
 G.1.2(4)
in Ada.Numerics.Generic_Elementary_Functions A.5.1(5)
 Tanh
in Ada.Numerics.Generic_Complex_Elementary_Functions
 G.1.2(6)
in Ada.Numerics.Generic_Elementary_Functions A.5.1(7)
 Time_Of

in Ada.Calendar 9.6(15)
in Ada.Calendar.Formatting 9.6.1(30/2), 9.6.1(31/2)
in Ada.Execution_Time D.14(9/2)
in Ada.Real_Time D.8(16)
Time_Of_Event
in Ada.Real_Time.Timing_Events D.15(6/2)
Time_Remaining
in Ada.Execution_Time.Timers D.14.1(8/2)
To_Ada
in Interfaces.C B.3(22), B.3(26), B.3(28), B.3(32), B.3(37),
 B.3(39), B.3(39.10/2), B.3(39.13/2), B.3(39.17/2),
 B.3(39.19/2), B.3(39.4/2), B.3(39.8/2)
in Interfaces.COBOL B.4(17), B.4(19)
in Interfaces.Fortran B.5(13), B.5(14), B.5(16)
To_Address
in System.Address_To_Access_Conversions 13.7.2(3)
in System.Storage_Elements 13.7.1(10)
To_Basic *in Ada.Characters.Handling* A.3.2(6), A.3.2(7)
To_Binary *in Interfaces.COBOL* B.4(45), B.4(48)
To_Bounded_String
in Ada.Strings.Bounded A.4.4(11)
To_C *in Interfaces.C* B.3(21), B.3(25), B.3(27), B.3(32),
 B.3(36), B.3(38), B.3(39.13/2), B.3(39.16/2), B.3(39.18/2),
 B.3(39.4/2), B.3(39.7/2), B.3(39.9/2)
To_Character
in Ada.Characters.Conversions A.3.4(5/2)
To_Chars_Ptr *in Interfaces.C.Strings* B.3.1(8)
To_COBOL *in Interfaces.COBOL* B.4(17), B.4(18)
To_Cursor *in Ada.Containers.Vectors* A.18.2(25/2)
To_Decimal *in Interfaces.COBOL* B.4(35), B.4(40), B.4(44),
 B.4(47)
To_Display *in Interfaces.COBOL* B.4(36)
To_Domain
in Ada.StringsMaps A.4.2(24)
in Ada.Strings.Wide_Maps A.4.7(24)
in Ada.Strings.Wide_Wide_Maps A.4.8(24/2)
To_Duration *in Ada.Real_Time* D.8(13)
To_Fortran *in Interfaces.Fortran* B.5(13), B.5(14), B.5(15)
To_Index *in Ada.Containers.Vectors* A.18.2(26/2)
To_Integer *in System.Storage_Elements* 13.7.1(10)
To_ISO_646 *in Ada.Characters.Handling* A.3.2(11), A.3.2(12)
To_Long_Binary *in Interfaces.COBOL* B.4(48)
To_Lower *in Ada.Characters.Handling* A.3.2(6), A.3.2(7)
To_Mapping
in Ada.StringsMaps A.4.2(23)
in Ada.Strings.Wide_Maps A.4.7(23)
in Ada.Strings.Wide_Wide_Maps A.4.8(23/2)
To_Packed *in Interfaces.COBOL* B.4(41)
To_Picture *in Ada.Text_IO.Editing* F.3.3(6)
To_Pointer
in System.Address_To_Access_Conversions 13.7.2(3)
To_Range
in Ada.StringsMaps A.4.2(24)
in Ada.Strings.Wide_Maps A.4.7(25)
in Ada.Strings.Wide_Wide_Maps A.4.8(25/2)
To_Ranges
in Ada.StringsMaps A.4.2(10)
in Ada.Strings.Wide_Maps A.4.7(10)
in Ada.Strings.Wide_Wide_Maps A.4.8(10/2)
To_Sequence
in Ada.StringsMaps A.4.2(19)
in Ada.Strings.Wide_Maps A.4.7(19)
in Ada.Strings.Wide_Wide_Maps A.4.8(19/2)
To_Set
in Ada.Containers.Hashed_Sets A.18.8(9/2)
in Ada.Containers.Ordered_Sets A.18.9(10/2)
in Ada.StringsMaps A.4.2(8), A.4.2(9), A.4.2(17),
 A.4.2(18)
in Ada.Strings.Wide_Maps A.4.7(8), A.4.7(9), A.4.7(17),
 A.4.7(18)
in Ada.Strings.Wide_Wide_Maps A.4.8(8/2), A.4.8(9/2),
 A.4.8(17/2), A.4.8(18/2)
To_String
in Ada.Characters.Conversions A.3.4(5/2)
in Ada.Strings.Bounded A.4.4(12)
in Ada.Strings.Unbounded A.4.5(11)
To_Time_Span *in Ada.Real_Time* D.8(13)
To_Unbounded_String
in Ada.Strings.Unbounded A.4.5(9), A.4.5(10)
To_Upper *in Ada.Characters.Handling* A.3.2(6), A.3.2(7)
To_Vector *in Ada.Containers.Vectors* A.18.2(13/2),
 A.18.2(14/2)
To_Wide_Character
in Ada.Characters.Conversions A.3.4(4/2), A.3.4(5/2)
To_Wide_String
in Ada.Characters.Conversions A.3.4(4/2), A.3.4(5/2)
To_Wide_Wide_Character
in Ada.Characters.Conversions A.3.4(4/2)
To_Wide_Wide_String
in Ada.Characters.Conversions A.3.4(4/2)
Translate
in Ada.Strings.Bounded A.4.4(53), A.4.4(54), A.4.4(55),
 A.4.4(56)
in Ada.Strings.Fixed A.4.3(18), A.4.3(19), A.4.3(20),
 A.4.3(21)
in Ada.Strings.Unbounded A.4.5(48), A.4.5(49), A.4.5(50),
 A.4.5(51)
Transpose
in Ada.Numerics.Generic_Complex_Arrays G.3.2(34/2)
in Ada.Numerics.Generic_Real_Arrays G.3.1(17/2)
Trim
in Ada.Strings.Bounded A.4.4(67), A.4.4(68), A.4.4(69)
in Ada.Strings.Fixed A.4.3(31), A.4.3(32), A.4.3(33),
 A.4.3(34)
in Ada.Strings.Unbounded A.4.5(61), A.4.5(62), A.4.5(63),
 A.4.5(64)
Unbounded_Slice
in Ada.Strings.Unbounded A.4.5(22.1/2), A.4.5(22.2/2)
Unchecked_Conversion
child of Ada 13.9(3)
Unchecked_Deallocation
child of Ada 13.11.2(3)
Union
in Ada.Containers.Hashed_Sets A.18.8(26/2), A.18.8(27/2)
in Ada.Containers.Ordered_Sets A.18.9(27/2), A.18.9(28/2)
Unit_Matrix
in Ada.Numerics.Generic_Complex_Arrays G.3.2(51/2)
in Ada.Numerics.Generic_Real_Arrays G.3.1(29/2)

Unit_Vector
in Ada.Numerics.Generic_Complex_Arrays G.3.2(24/2)
in Ada.Numerics.Generic_Real_Arrays G.3.1(14/2)

Update *in Interfaces.C.Strings* B.3.1(18), B.3.1(19)

Update_Element
in Ada.Containers.Doubly_Linked_Lists A.18.3(17/2)
in Ada.Containers.Hashed_Maps A.18.5(17/2)
in Ada.Containers.Ordered_Maps A.18.6(16/2)
in Ada.Containers.Vectors A.18.2(33/2), A.18.2(34/2)

Update_Element_Preserving_Key
in Ada.Containers.Hashed_Sets A.18.8(58/2)
in Ada.Containers.Ordered_Sets A.18.9(73/2)

Update_Error *in Interfaces.C.Strings* B.3.1(20)

UTC_Time_Offset
in Ada.Calendar.Time_Zones 9.6.1(6/2)

Valid
in Ada.Text_IO.Editing F.3.3(5), F.3.3(12)
in Interfaces.COBOL B.4(33), B.4(38), B.4(43)

Value
in Ada.Calendar.Formatting 9.6.1(36/2), 9.6.1(38/2)
in Ada.Environment_Variables A.17(4/2)
in Ada.Numerics.Discrete_Random A.5.2(26)
in Ada.Numerics.Float_Random A.5.2(14)
in Ada.Strings.Maps A.4.2(21)
in Ada.Strings.Wide_Maps A.4.7(21)
in Ada.Strings.Wide_Wide_Maps A.4.8(21/2)
in Ada.Task_Attributes C.7.2(4)
in Interfaces.C.Pointers B.3.2(6), B.3.2(7)

in Interfaces.C.Strings B.3.1(13), B.3.1(14), B.3.1(15),
B.3.1(16)

Virtual_Length
in Interfaces.C.Pointers B.3.2(13)

Wide_Hash
child of Ada.Strings.Wide_Bounded A.4.7(1/2)
child of Ada.Strings.Wide_Fixed A.4.7(1/2)
child of Ada.Strings.Wide_Unbounded A.4.7(1/2)

Wide_Exception_Name *in Ada.Exceptions* 11.4.1(2/2),
11.4.1(5/2)

Wide_Expanded_Name *in Ada.Tags* 3.9(7/2)

Wide_Wide_Hash
child of Ada.Strings.Wide_Wide_Bounded A.4.8(1/2)
child of Ada.Strings.Wide_Wide_Fixed A.4.8(1/2)
child of Ada.Strings.Wide_Wide_Unbounded A.4.8(1/2)

Wide_Wide_Exception_Name
in Ada.Exceptions 11.4.1(2/2), 11.4.1(5/2)

Wide_Wide_Expanded_Name *in Ada.Tags* 3.9(7/2)

Write
in Ada.Direct_IO A.8.4(13)
in Ada.Sequential_IO A.8.1(12)
in Ada.Storage_IO A.9(7)
in Ada.Streams 13.13.1(6)
in Ada.Streams.Stream_IO A.12.1(18), A.12.1(19)
in System.RPC E.5(8)

Year
in Ada.Calendar 9.6(13)
in Ada.Calendar.Formatting 9.6.1(21/2)

Q.4 Language-Defined Exceptions

This clause lists all language-defined exceptions.

1/2

Argument_Error
in Ada.Numerics A.5(3/2)

Communication_Error
in System.RPC E.5(5)

Constraint_Error
in Standard A.1(46)

Conversion_Error
in Interfaces.COBOL B.4(30)

Data_Error
in Ada.Direct_IO A.8.4(18)
in Ada.IO_Exceptions A.13(4)
in Ada.Sequential_IO A.8.1(15)
in Ada.Storage_IO A.9(9)
in Ada.Streams.Stream_IO A.12.1(26)
in Ada.Text_IO A.10.1(85)

Device_Error
in Ada.Direct_IO A.8.4(18)
in Ada.Directories A.16(43/2)
in Ada.IO_Exceptions A.13(4)
in Ada.Sequential_IO A.8.1(15)
in Ada.Streams.Stream_IO A.12.1(26)
in Ada.Text_IO A.10.1(85)

Dispatching_Policy_Error
in Ada.Dispatching D.2.1(1.2/2)

End_Error
in Ada.Direct_IO A.8.4(18)
in Ada.IO_Exceptions A.13(4)
in Ada.Sequential_IO A.8.1(15)
in Ada.Streams.Stream_IO A.12.1(26)
in Ada.Text_IO A.10.1(85)

Group_Budget_Error
in Ada.Execution_Time.Group_Budgets D.14.2(11/2)

Index_Error
in Ada.Strings A.4.1(5)

Layout_Error
in Ada.IO_Exceptions A.13(4)
in Ada.Text_IO A.10.1(85)

Length_Error
in Ada.Strings A.4.1(5)

Mode_Error
in Ada.Direct_IO A.8.4(18)
in Ada.IO_Exceptions A.13(4)
in Ada.Sequential_IO A.8.1(15)
in Ada.Streams.Stream_IO A.12.1(26)
in Ada.Text_IO A.10.1(85)

Name_Error	Storage_Error
<i>in Ada.Direct_IO</i> A.8.4(18)	<i>in Standard</i> A.1(46)
<i>in Ada.Directories</i> A.16(43/2)	Tag_Error
<i>in Ada.IO_Exceptions</i> A.13(4)	<i>in Ada.Tags</i> 3.9(8)
<i>in Ada.Sequential_IO</i> A.8.1(15)	Tasking_Error
<i>in Ada.Streams.Stream_IO</i> A.12.1(26)	<i>in Standard</i> A.1(46)
<i>in Ada.Text_IO</i> A.10.1(85)	Terminator_Error
Pattern_Error	<i>in Interfaces.C</i> B.3(40)
<i>in Ada.Strings</i> A.4.1(5)	Time_Error
Picture_Error	<i>in Ada.Calendar</i> 9.6(18)
<i>in Ada.Text_IO.Editing</i> F.3.3(9)	Timer_Resource_Error
Pointer_Error	<i>in Ada.Execution_Time.Timers</i> D.14.1(9/2)
<i>in Interfaces.C.Pointers</i> B.3.2(8)	Translation_Error
Program_Error	<i>in Ada.Strings</i> A.4.1(5)
<i>in Standard</i> A.1(46)	Unknown_Zone_Error
Status_Error	<i>in Ada.Calendar.Time_Zones</i> 9.6.1(5/2)
<i>in Ada.Direct_IO</i> A.8.4(18)	Use_Error
<i>in Ada.Directories</i> A.16(43/2)	<i>in Ada.Direct_IO</i> A.8.4(18)
<i>in Ada.IO_Exceptions</i> A.13(4)	<i>in Ada.Directories</i> A.16(43/2)
<i>in Ada.Sequential_IO</i> A.8.1(15)	<i>in Ada.IO_Exceptions</i> A.13(4)
<i>in Ada.Streams.Stream_IO</i> A.12.1(26)	<i>in Ada.Sequential_IO</i> A.8.1(15)
<i>in Ada.Text_IO</i> A.10.1(85)	<i>in Ada.Streams.Stream_IO</i> A.12.1(26)
	<i>in Ada.Text_IO</i> A.10.1(85)

Q.5 Language-Defined Objects

1/2 This clause lists all language-defined constants, variables, named numbers, and enumeration literals.

ACK <i>in Ada.Characters.Latin_1</i> A.3.3(5)
Acute <i>in Ada.Characters.Latin_1</i> A.3.3(22)
Ada_To_COBOL <i>in Interfaces.COBOL</i> B.4(14)
Alphanumeric_Set
<i>in Ada.Strings.Maps.Constants</i> A.4.6(4)
Ampersand <i>in Ada.Characters.Latin_1</i> A.3.3(8)
APC <i>in Ada.Characters.Latin_1</i> A.3.3(19)
Apostrophe <i>in Ada.Characters.Latin_1</i> A.3.3(8)
Asterisk <i>in Ada.Characters.Latin_1</i> A.3.3(8)
Basic_Map
<i>in Ada.Strings.Maps.Constants</i> A.4.6(5)
Basic_Set
<i>in Ada.Strings.Maps.Constants</i> A.4.6(4)
BEL <i>in Ada.Characters.Latin_1</i> A.3.3(5)
BPH <i>in Ada.Characters.Latin_1</i> A.3.3(17)
Broken_Bar <i>in Ada.Characters.Latin_1</i> A.3.3(21)
BS <i>in Ada.Characters.Latin_1</i> A.3.3(5)
Buffer_Size <i>in Ada.Storage_IO</i> A.9(4)
CAN <i>in Ada.Characters.Latin_1</i> A.3.3(6)
CCH <i>in Ada.Characters.Latin_1</i> A.3.3(18)
Cedilla <i>in Ada.Characters.Latin_1</i> A.3.3(22)
Cent_Sign <i>in Ada.Characters.Latin_1</i> A.3.3(21)
char16_nul <i>in Interfaces.C</i> B.3(39.3/2)
char32_nul <i>in Interfaces.C</i> B.3(39.12/2)
CHAR_BIT <i>in Interfaces.C</i> B.3(6)
Character_Set
<i>in Ada.Strings.Wide_Maps</i> A.4.7(46/2)
<i>in Ada.Strings.Wide_Maps.Wide_Constants</i> A.4.8(48/2)

Circumflex <i>in Ada.Characters.Latin_1</i> A.3.3(12)
COBOL_To_Ada <i>in Interfaces.COBOL</i> B.4(15)
Colon <i>in Ada.Characters.Latin_1</i> A.3.3(10)
Comma <i>in Ada.Characters.Latin_1</i> A.3.3(8)
Commercial_At
<i>in Ada.Characters.Latin_1</i> A.3.3(10)
Control_Set
<i>in Ada.Strings.Maps.Constants</i> A.4.6(4)
Copyright_Sign
<i>in Ada.Characters.Latin_1</i> A.3.3(21)
CPU_Tick <i>in Ada.Execution_Time</i> D.14(4/2)
CPU_Time_First <i>in Ada.Execution_Time</i> D.14(4/2)
CPU_Time_Last <i>in Ada.Execution_Time</i> D.14(4/2)
CPU_Time_Unit <i>in Ada.Execution_Time</i> D.14(4/2)
CR <i>in Ada.Characters.Latin_1</i> A.3.3(5)
CSI <i>in Ada.Characters.Latin_1</i> A.3.3(19)
Currency_Sign
<i>in Ada.Characters.Latin_1</i> A.3.3(21)
DC1 <i>in Ada.Characters.Latin_1</i> A.3.3(6)
DC2 <i>in Ada.Characters.Latin_1</i> A.3.3(6)
DC3 <i>in Ada.Characters.Latin_1</i> A.3.3(6)
DC4 <i>in Ada.Characters.Latin_1</i> A.3.3(6)
DCS <i>in Ada.Characters.Latin_1</i> A.3.3(18)
Decimal_Digit_Set
<i>in Ada.Strings.Maps.Constants</i> A.4.6(4)
Default_Aft
<i>in Ada.Text_IO</i> A.10.1(64), A.10.1(69), A.10.1(74)
<i>in Ada.Text_IO.Complex_IO</i> G.1.3(5)

Default_Base *in Ada.Text_IO* A.10.1(53), A.10.1(58)
 Default_Bit_Order *in System* 13.7(15/2)
 Default_Currency
in Ada.Text_IO.Editing F.3.3(10)
 Default_Deadline
in Ada.Dispatching.EDF D.2.6(9/2)
 Default_Exp
in Ada.Text_IO A.10.1(64), A.10.1(69), A.10.1(74)
in Ada.Text_IO.Complex_IO G.1.3(5)
 Default_Fill *in Ada.Text_IO.Editing* F.3.3(10)
 Default_Fore
in Ada.Text_IO A.10.1(64), A.10.1(69), A.10.1(74)
in Ada.Text_IO.Complex_IO G.1.3(5)
 Default_Priority *in System* 13.7(17)
 Default_Quantum
in Ada.Dispatching.Round_Robin D.2.5(4/2)
 Default_Radix_Mark
in Ada.Text_IO.Editing F.3.3(10)
 Default_Separator
in Ada.Text_IO.Editing F.3.3(10)
 Default_Setting *in Ada.Text_IO* A.10.1(80)
 Default_Width *in Ada.Text_IO* A.10.1(53), A.10.1(58),
 A.10.1(80)
 Degree_Sign *in Ada.Characters.Latin_1* A.3.3(22)
 DEL *in Ada.Characters.Latin_1* A.3.3(14)
 Diaeresis *in Ada.Characters.Latin_1* A.3.3(21)
 Division_Sign
in Ada.Characters.Latin_1 A.3.3(26)
 DLE *in Ada.Characters.Latin_1* A.3.3(6)
 Dollar_Sign *in Ada.Characters.Latin_1* A.3.3(8)
 e *in Ada.Numerics* A.5(3/2)
 EM *in Ada.Characters.Latin_1* A.3.3(6)
 Empty_List
in Ada.Containers.Doubly_Linked_Lists A.18.3(8/2)
 Empty_Map
in Ada.Containers.Hashed_Maps A.18.5(5/2)
in Ada.Containers.Ordered_Maps A.18.6(6/2)
 Empty_Set
in Ada.Containers.Hashed_Sets A.18.8(5/2)
in Ada.Containers.Ordered_Sets A.18.9(6/2)
 Empty_Vector
in Ada.Containers.Vectors A.18.2(10/2)
 ENQ *in Ada.Characters.Latin_1* A.3.3(5)
 EOT *in Ada.Characters.Latin_1* A.3.3(5)
 EPA *in Ada.Characters.Latin_1* A.3.3(18)
 Equals_Sign *in Ada.Characters.Latin_1* A.3.3(10)
 ESA *in Ada.Characters.Latin_1* A.3.3(17)
 ESC *in Ada.Characters.Latin_1* A.3.3(6)
 ETB *in Ada.Characters.Latin_1* A.3.3(6)
 ETX *in Ada.Characters.Latin_1* A.3.3(5)
 Exclamation *in Ada.Characters.Latin_1* A.3.3(8)
 Failure *in Ada.Command_Line* A.15(8)
 Feminine_Ordinal_Indicator
in Ada.Characters.Latin_1 A.3.3(21)
 FF *in Ada.Characters.Latin_1* A.3.3(5)
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