Student Number

SOLUTION

The University of Melbourne

Department of Computing and Information Systems

CRYPTOGRAPHY AND SECURITY

September, 2017

Quiz Duration: 45 minutes.

Length: This paper has 7 pages including this cover page.

Authorised Materials: None.

Instructions to Students: Answer all questions in this exam booklet. Total marks for the test is 50. This is worth 10% of the final mark in the subject;

Calculators: No Calculators are permitted.

Library: This paper must be returned and not taken out of the exam hall.

- 1. (10 marks) Short Answer Questions (Please answer in the space provided).
 - (a) Let p be a prime number. Then for any x, $x^p \mod p = \dots \times \mod p$
 - (b) (15 19) mod 26= **22**

 - (d) $2^{30}6^{2600}5^{33} \mod 7 = -1$ or 6
 - (e) $2^{144}3^{132}5^{100} \mod 4 = 0$
 - (f) $\phi(p_1 \ p_2) = (p_1 1)(p_2 1)$ where p_1 and p_2 are distinct primes and ϕ is Euler's function.
 - (g) For any positive integer k, $\phi(p^k) = P^{k-1}(p-1)$ where p is a prime number and ϕ is Euler's function.
 - (h) The minimum positive integer x that satisfies the following relations is $x = 2 \mod 7$; $x = 3 \mod 5$.
 - (i) All encryption algorithms are based on two general principles: substitution and permutation / transposition
 - (j) An...activeattack attempts to alter system resources or affect their operation.

- 2. (8 marks) RSA and Public Key Crypto systems.
 - (a) What are the hard mathematical problems on which security of RSA cryptosystem is based? You need to define the problems, not just the names.
 - O integer factorisation

 hard to factorise a large integer into product
 of primes
 - ② RSA problem
 given C and e, hard to find M
 where C = Me mod n
 - (b) What are the hard mathematical problems on which security of Diffie-Hellman Key Agreement protocol is based? You need to define the problems, not just the names.
 - ① discrete logarithm

 let $g^* = y$, it is hard to solve x Knowing g and y
 - © computational Diffie-Hellman problem given g, ga, gb it is hard to find gab

Alice wants to configure her RSA parameters. She chooses two large random primes p and q. Fill in the blanks in the following items which will help her compute the RSA parameters.

- i. Alice's RSA modulus n is -
- ii. The encryption exponent e is chosen such that $\gcd(e, \varphi(n)) = 1$
- iii. The decryption exponent d is found such that $ed = 1 \mod Q(n)$
- iv. The ciphertext for the message m is _____ mod _____

- 3. (12 marks) This question is about computing the inverse of a number modulo n, where n a positive integer. Note: Inverse of a number $a \mod n$ is a number $a \mod n$ such that $a = 1 \mod n$.
 - (a) The Extended GCD algorithm (XGCD), also known as the Euclidean algorithm, takes two given integers a and b as inputs and returns three integers g, x and y such that

$$a x + b y = g$$
,

where g is the greatest common divisor of the input integers. Write a pseudocode for the function **inverse modulo** n using the XGCD function given above. NOTE: There is no need for you write XGCD function.

- (b) You have been given the results from the XGCD function below:
 - i. XGCD(12987, 46799) = 1, -13488, 3743
 - ii. XGCD(12, 39) = 3, -3, 1
 - iii. XGCD(17, 29) = 1, 12, -7

Now determine the inverse of the following numbers:

- i. 12 mod 39 N/A
- ii. 12987 mod 46799 13488 or 33311
- iii. 17 mod 29 12
- iv. 12 mod 17

4. (10 marks) For the prime numbers p=11 and q=7, calculate the non-trivial RSA keys e and d, e>1, satisfying the condition that d has the smallest possible values.

$$\begin{aligned}
N &= PQ = 77 \\
Q(n) &= Q(P) Q(Q) = 10 \times 6 = 60 \\
d &= 7 \\
60 &= 8 \times 7 + 4 \\
7 &= 1 \times 4 + 3 \\
4 &= 1 \times 3 + 1
\end{aligned}$$

$$\begin{aligned}
1 &= 4 - 1 \times 3 \\
1 &= 4 - 1 \times (7 - 1 \times 4) \\
1 &= 2 \times 4 - 1 \times 7 \\
1 &= 2 \times (60 - 8 \times 7) - 1 \times 7 \\
1 &= 2 \times 60 - 17 \times 7
\end{aligned}$$

$$e &= + 43$$

5. (10 marks) The following equations and figure describe one of the standard modes of usage of symmetric key encryption.

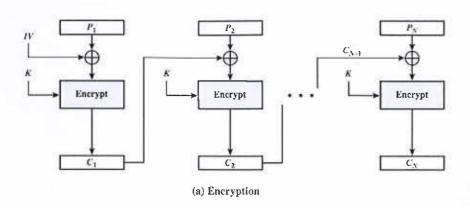


Figure 1: A Standard Mode of Encryption

Encryption:

$$C_1 = (E_K[IV \oplus P_1]).$$

$$C_j = (E_K[C_{j-1} \oplus P_j]), j > 1.$$

(a) What is the name of this mode?

Cipher Block Chaining

(b) Expand the abbreviations and functions used in the equations:

i.
$$IV =$$
 initialisation vector
ii. $K =$ secret key
iii. $E_y[x] =$ encryption function on x with key y

(c) Complete the equations for decryption below:

Decryption:
$$P_{1} = \underbrace{\qquad \qquad}_{P_{j} = ---} D_{k} [C_{1}] \oplus IV$$

$$P_{j} = \underbrace{\qquad \qquad}_{P_{k} [C_{j}]} \oplus C_{j-1} \qquad j \gg 2$$
(d) What is the effect on the plain text of a one bit error in the transmission of an engage of the effect of the plain text of a one bit error in the transmission of an engage of the effect of the plain text of a one bit error in the transmission of an engage of the effect of the effect of the plain text of a one bit error in the transmission of the effect of t

of an encrypted "block C_i "?

END OF EXAMINATION