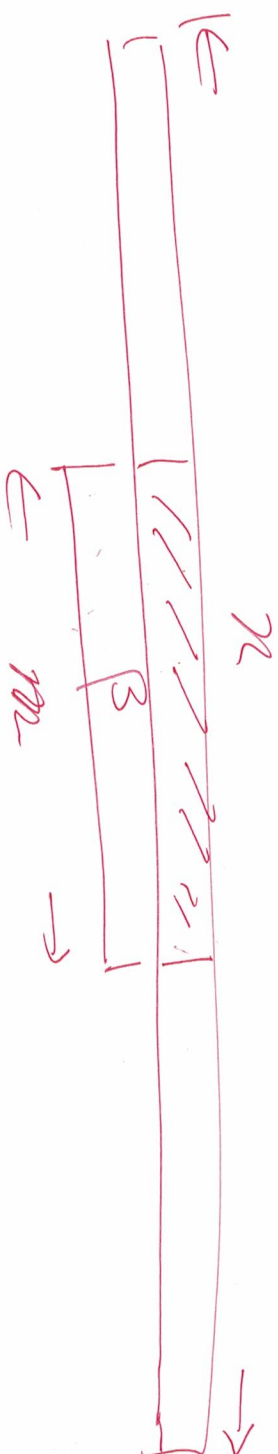


Cpts 515. 10/30/2020.

Last time: Karp-Rabin Alg.



$$h(\underbrace{a}_{m+i-1}) = h(b)$$

ε : collision probability. We control the

$$\text{total error probability } \varepsilon \leq \frac{1}{2} \dots (*)$$

In above, a and b are m -bit unsigned numbers.

$$h(a) = h(b) \Rightarrow \underbrace{a-b}_{*} \bmod q = 0$$

all possible values of $|a-b|$.

↳ how many: 2^m values.

So, $|a-b|$ has at most m prime divisors.
how many g 's to make (*) true? at m many.

We need find N s.t. $O(1) N$ contains

2^{nm} primes. So that the error \leq will be

$$\leq \frac{m}{2^{nm}} = \frac{1}{2^n} \text{ to make } (**) \text{ true.}$$

How big is N ?

Define $\pi(N)$ = "the # of primes $\leq N$ ".

From Euler, $\pi(N) \sim \frac{N}{\ln N}$.

Better result (Rosser & Schoenfeld): when $N \geq 17$,

$$\frac{N}{\ln N} \leq \pi(N) \leq 1.25506 \frac{N}{\ln N}.$$

Now, $2nm \leq \frac{N}{\ln N}$.

$$\Rightarrow N = (20nm) \ln(20nm).$$

↳ This is small!

Mid-term

Linear algebra.

Linear transform! but this is not
the most useful topic for CS.

Matrix = Graph

Graph: ① path count (up to ∞ length)

② "density" is a graph

③ Graph similarity / graph isomorphism.

...

Let G be a graph (assuming that G is connected) with n nodes. M is G 's adjacency matrix:

$$M[i, j] = 1 \text{ if node } i \rightarrow \text{node } j \text{ is an edge,}$$

$$M[i, j] = 0 \text{ if } \dots \dots \dots \text{is not an edge.}$$

①. M has a unique and largest eigenvalue, called Perron number, denoted by λ , (a positive real #).

② All other eigenvalues of M are $< \lambda$, (i.e., $|\lambda'| < \lambda$ for all other eigenvalues λ').
L_{norm} of λ' .

A tool in your toolbox:

A is $m \times m$ matrix. A is positive. (≥ 0)

Eigenvalue λ : $AX = \lambda X$.

you have n eigenvalues.

left eigenvector $uA = \lambda u$.

right eigenvector $Av = \lambda v$.

Matlab can take care of these.

Next,

$$① \quad M^n =_{\text{adj}} \underbrace{M \cdot M \cdot M \cdot \dots \cdot M}_n \text{ times.}$$

$$② \quad M^n[i, j] = \text{the total \# of walks from node } i \text{ to node } j \text{ in } G,$$

$$③ \quad M^n \text{ can be approx. by}$$

$$\begin{array}{c} \lambda^n \\ \swarrow \quad \searrow \\ V \quad U^T \\ \text{Perron vector} \quad \text{left eigenvector of } \lambda \end{array} \quad \xrightarrow{\text{Transpose}} \quad \begin{array}{c} \text{right eigenvector of } \lambda \end{array}$$

④ The total # of calls from node i to node j with length n , $M^n \Sigma_{i,j}$, can be approx. by

$$\frac{v_i \cdot u_j}{\|u\|} \cdot \lambda^n \cdot v u^T$$

where $\|u\| = \sum_k u_k$, and v_i, u_j are the components in vectors v, u .

⑤ The total # of calls from node i with length n , can be approx. by

$$\frac{2k_i}{\|u\|} \cdot \lambda^n \cdot v u^T \cdot \dots \cdot (*)$$

Now

$$\limsup_{n \rightarrow \infty} \frac{\log |S_n|}{n} = \log \lambda \quad \text{using } (*)$$

↓
Perron number

where S_n : The total # of walks for N_0 .

What if G has multiple SCC?

Each SCC has a λ .

take max of all the λ 's.

$$\limsup_{n \rightarrow \infty} \frac{\log |S_n|}{n}$$

$$= \log \lambda$$

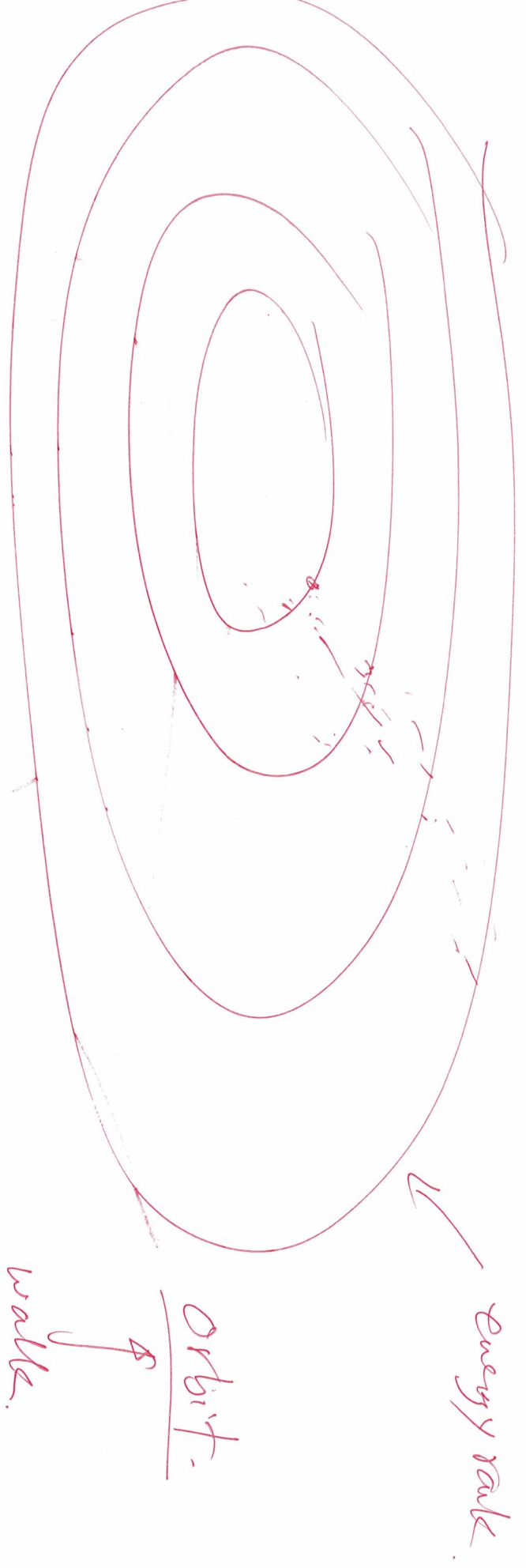
↑
we know how to get this.

For different applications, you need interpret the S_n

differently:

①. S_n is the # of walks with $|E_n| = n$.

②. "walk" ————— example.



Energy: high \Rightarrow more ways to walk
(# of orbits is high).