

CPTS 515. 10/12/2020.

Today: Automata Approach.

Background: Undergraduate CPTS 317.
(Automata & Formal Languages).

1.0. Automata: Basic Concept.

Automaton: anything that's moving.
(and not moving).


So, it's universal.
↳ dead automaton.

History: 1st Automaton = Turing machine.

However, TMs are too powerful.
↳ Halting is undecidable.

①, In reality, we need weak model so that many good ~~properties~~^{properties} can be decided using Algorithms;

②, We need study a hierarchy of such models,

- Finite Automata  most useful,
Pushdown Automata

⋮

③. Finite Automata can be interpreted in many ways.

Any C-program that fixed and finite amount of memory is a finite automata.



Example:

int x, y;

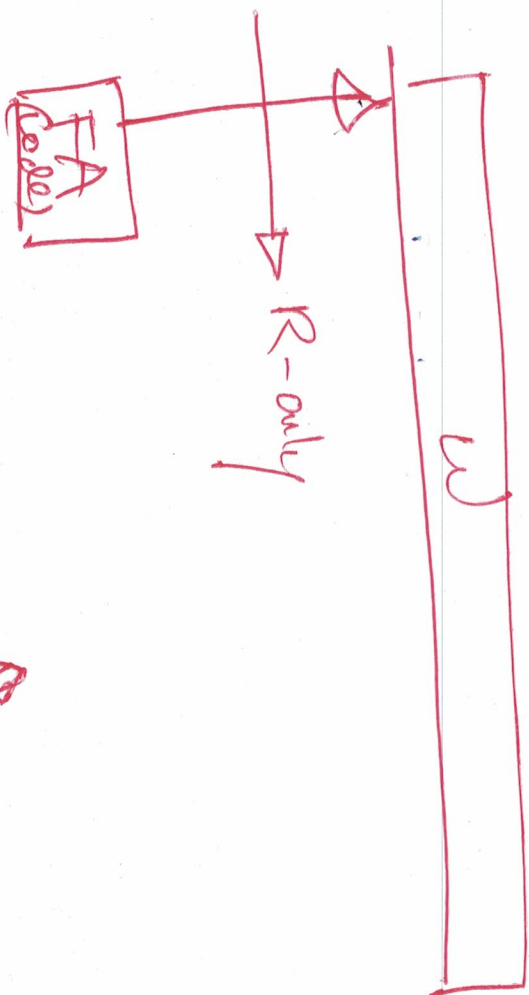
unbounded memory.

...

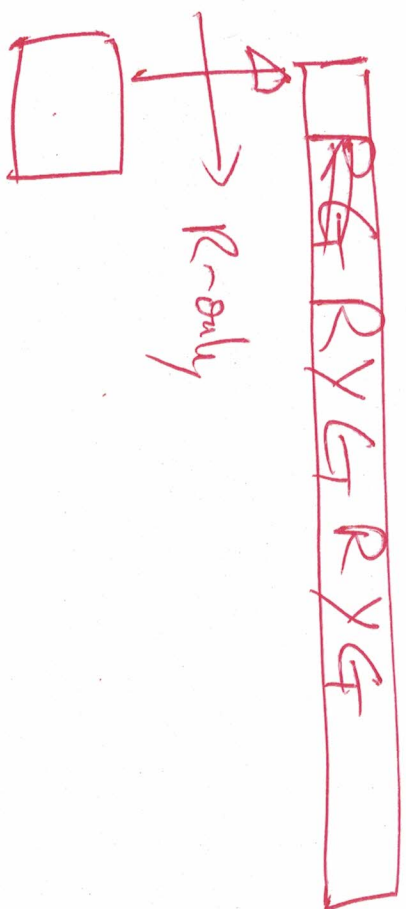
bool x, y;

finite or bounded memory

One way to avoid a FA:



Example: describe a FA to accept



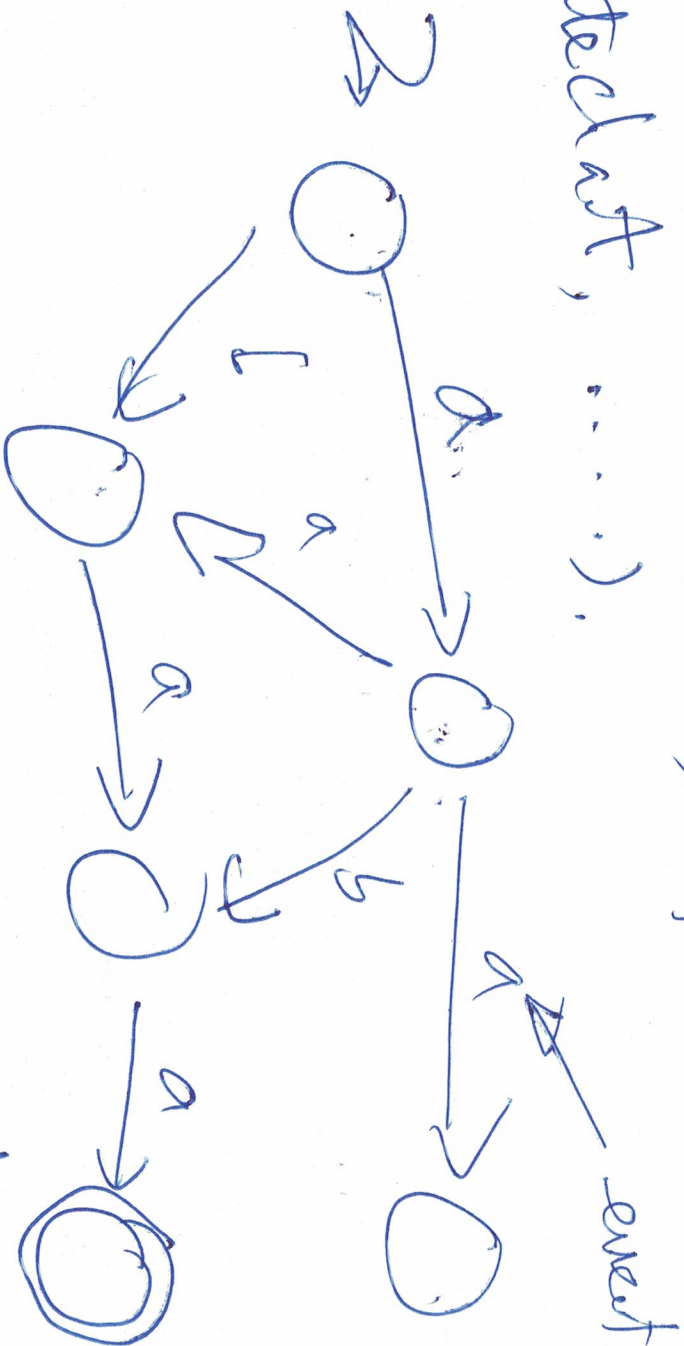
(R Y G)*

First I need R,
then I need Y,
then I need G,
Repeat ...

how much
memory?

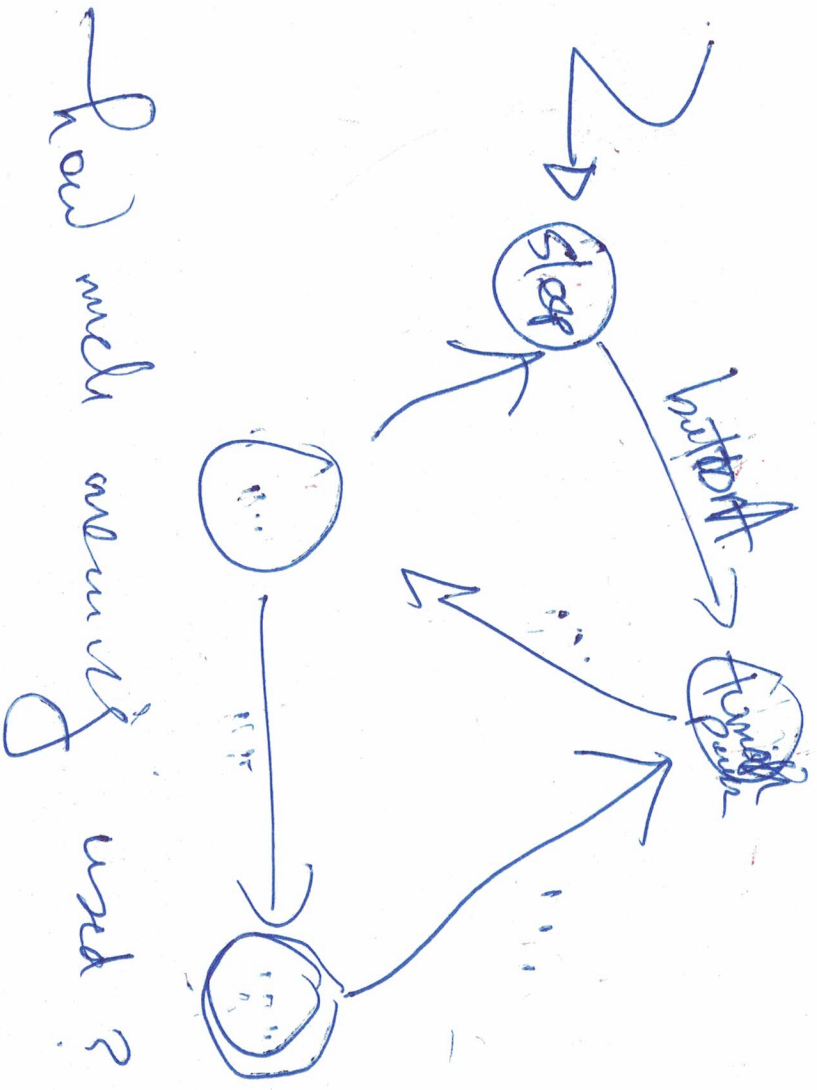
The money used here is to tell one of three colors. $\log_2 3$ bits.

Another way to describe FA: a graph representation (state-transition diagram; finite state transition system) (statechart, ...).



two versions; det. vs. nondet.

(A good name shall be: labeled transition system, LTS).



after member,

which?
current
state.

how much memory used?

$\log_2 4$

bits.

why? you only need

Events don't take memory!

to remember one of 4 states,
at any time.

Go back to SCC:
↳ loop-analysis of a graph

logic thunks:

C-programs (with fite memory) \rightarrow fite instructions \rightarrow
labeled transition \rightarrow a graph.

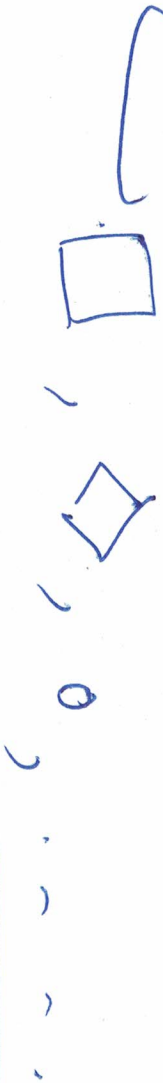
From this chain:

if you want an alg to solve problems
on C-programs, we have a new

approach now:

(if it runs on fite memory) translate
the program into a graph and use
known graph alg.

(Vardi & LTL): uses SCC to solve a program's liveness query:

□, ◇, ○, ...

Next time: ideas in using negation in automata.