Lecture 3: Rademacher Complexity and VC-Dimension

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Dec. 18th., 2022

Contents

1	PAG	C Learnability for Infinite Hypothese Set	2	
2	Rad	lemacher complexity	lexity 2	
3	VC			
	3.1	Definition of VC-Dimension	5	
	3.2	Growth Function	6	
	3.3	Relate Growth Function to Radmatcher Complexity	8	
	3.4	Generalization Bounds via Growth Function and VC-Dimension	9	
	3.5	Examples	9	
	3.6	Lower Bounds	9	

1 PAC Learnability for Infinite Hypothese Set

• Remark (Universal Consistency for Infinite Hypothesis Set)

When $|\mathcal{H}| < \infty$, we can use sample complexity bounds that involve $\log |\mathcal{H}|$ for universal consistency of ERM. Obviously, we cannot do so when $|\mathcal{H}| = \infty$. A general idea of analyzing infinite hypothesis set consists of reducing the infinite case to the analysis of finite sets of hypotheses and then proceed as in the previous chapter.

There are different techniques for that reduction, each relying on a different notion of complexity for the family of hypotheses.

2 Rademacher complexity

• Remark (Notations)

We will continue to use \mathcal{H} to denote a *hypothesis set* as in the previous chapters, and $h \in \mathcal{H}$ an *element* of \mathcal{H} . Many of the results of this section are general and hold for an arbitrary loss function $L: \mathcal{Y} \times \mathcal{Y} \to \mathbb{R}$. To each $h: \mathcal{X} \to \mathcal{Y}$, we can associate a function g:

$$g:(x,y)\in\mathcal{X}\times\mathcal{Y}\to L(h(x),y)$$

without explicitly describing the specific loss L used. In what follows \mathcal{G} will generally be interpreted as **the family of loss functions** associated to \mathcal{H} .

• Definition (Empirical Rademacher Complexity)

Let \mathcal{G} be a family of functions mapping from \mathcal{Z} to [a,b] and $\mathcal{D}_m = (z_1,\ldots,z_m)$ a fixed sample of size m with elements in \mathcal{Z} . Then, the empirical Rademacher complexity of \mathcal{G} with respect to the sample \mathcal{D} is defined as:

$$\widehat{\mathfrak{R}}_{\mathcal{D}}(\mathcal{G}) = \mathbb{E}_{\sigma} \left[\sup_{g \in \mathcal{G}} \frac{1}{m} \sum_{i=1}^{m} \sigma_{i} g(z_{i}) \right]$$
(1)

where $\sigma := (\sigma_1, \dots, \sigma_m)$, with σ_i s independent uniform random variables taking values in $\{-1, +1\}$. The random variables σ_i are called <u>Rademacher variables</u>.

• Remark (How Well to Fit Random Noise)

The Rademacher complexity captures the richness of a family of functions by measuring the degree to which a hypothesis set can fit random noise. The richer or more complex families \mathcal{G} can generate more vectors $g_{\mathcal{D}}$ and thus better correlate with random noise, on \mathcal{D}_m .

The intuition is that if a hypothesis set can fit arbitrary noise, then it is too large to bound the performance of ERM, i.e. it is very likely to have overfitting (zero empirical error but arbitrary bad generalization error).

• Remark

$$\widehat{\mathfrak{R}}_{\mathcal{D}}(\mathcal{G}) = \mathbb{E}_{\sigma} \left[\sup_{g \in \mathcal{G}} \frac{\langle \sigma_m, g_{\mathcal{D}_m} \rangle}{m} \right]$$

which measures the correlation between the random noise $\sigma_m := \{\sigma_i\}$ and $g_{\mathcal{D}_m} := \{g(Z_i)\}_{i=1}^m$. The supremum $\sup_{g \in \mathcal{G}} \langle \sigma_m, g_{\mathcal{D}_m} \rangle / m$ is a measure of how well the function class \mathcal{G} correlates

with σ over the sample \mathcal{D}_m . Thus, the empirical Rademacher complexity measures on average how well the function class \mathcal{G} correlates with random noise on \mathcal{D}_m .

• Remark (Rademacher Variables = Symmetric Bernoulli Random Variable) A Rademacher variables is a random variable on $\{-1,+1\}$ with symmetric probability $P(\sigma_i = +1) = P(\sigma_i = -1) = \frac{1}{2}$, i.e. a set of i.i.d. symmetric Bernoulli random variables.

A Rademacher process is a stochastic process $\{\sigma_i g(X_i)\}$ conditioning on $\{X_i\}$, which is subgaussian process.

• Definition (Rademacher Complexity)

Let $\mathcal{P}_{X,Y}$ denote the distribution according to which samples $\mathcal{D}_m = \{(X_i, Y_i), 1 \leq i \leq m\}$ are drawn, $\mathcal{H} \subset \mathcal{C}$ be a hypthesis class where $h \in \mathcal{H}$ is a function from \mathcal{X} to \mathcal{Y} , \mathcal{G} be a family of loss functions associated to \mathcal{H} .

For any integer $m \geq 1$, the Rademacher complexity of \mathcal{G} is defined as the expectation of the empirical Rademacher complexity over all samples of size m drawn according to $\mathcal{P}_{X,Y}$:

$$\mathfrak{R}_m(\mathcal{G}) = \mathbb{E}_{\mathcal{P}_{X,Y}^m} \left[\widehat{\mathfrak{R}}_{\mathcal{D}_m}(\mathcal{G}) \right].$$

• Proposition 2.1 (Consistency Bound with respect to Rademacher Complexity)
[Mohri et al., 2018]

Let \mathcal{G} be a family of functions mapping from \mathcal{Z} to [0,1]. Then, for any $\delta > 0$, with probability at least $1 - \delta$, each of the following holds for all $g \in \mathcal{G}$:

$$\mathbb{E}\left[g(Z)\right] \le \frac{1}{m} \sum_{i=1}^{m} g(Z_i) + 2\mathfrak{R}_m(\mathcal{G}) + \sqrt{\frac{\log(1/\delta)}{2m}}$$
 (2)

and

$$\mathbb{E}\left[g(Z)\right] \le \frac{1}{m} \sum_{i=1}^{m} g(Z_i) + 2\widehat{\mathfrak{R}}_{\mathcal{D}_m}(\mathcal{G}) + 3\sqrt{\frac{\log(2/\delta)}{2m}}$$
(3)

• Remark Define a function Φ on \mathcal{D}_m by

$$\Phi(\mathcal{D}_m) := \sup_{g \in \mathcal{G}} \left\{ \mathbb{E}\left[g(Z)\right] - \widehat{\mathbb{E}}_{\mathcal{D}_m}\left[g(Z)\right] \right\},\,$$

where $\widehat{\mathbb{E}}_{\mathcal{D}_m}[g(Z)] = \frac{1}{m} \sum_{i=1}^m g(Z_i)$. This function has bounded difference if only one sample makes changes. The proof consists of three parts:

1. Bound the probability of tail event of Φ

$$\Phi(\mathcal{D}_m) - \mathbb{E}_{\mathcal{P}_Z^m} \left[\Phi(\mathcal{D}_m)\right]$$

2. Bound the expectation $\mathbb{E}_{\mathcal{P}_{Z}^{m}}[\Phi(\mathcal{D}_{m})]$ by Radmatcher complexity on \mathcal{G}

$$\mathbb{E}_{\mathcal{P}_Z^m}\left[\Phi(\mathcal{D}_m)\right] \leq 2\mathfrak{R}_m(\mathcal{G}).$$

This obtain the inequality (2).

3. Bound the probability of tail events for empircal Radmatcher complexity:

$$\mathfrak{R}_m(\mathcal{G}) - \widehat{\mathfrak{R}}_{\mathcal{D}_m}(\mathcal{G})$$

Combines the difference in (1) and (3) using union bounds.

• Lemma 2.2 (Radmatcher Complexity for 0-1 Loss) [Mohri et al., 2018] Let \mathcal{H} be a family of functions taking values in $\{-1,+1\}$ and let \mathcal{G} be the family of loss functions associated to \mathcal{H} for the zero-one loss:

$$\mathcal{G} := \left\{ (x, y) \mapsto \mathbb{1}_{h(x) \neq y} : h \in \mathcal{H} \right\}.$$

For any sample $\mathcal{D} = ((X_1, Y_1), \dots, (X_m, Y_m))$ of elements in $\mathcal{X} \times \{-1, +1\}$, let $\mathcal{D}_{\mathcal{X}} := (X_1, \dots, X_m)$. Then, the following relation holds between **the empirical Rademacher** complexities of \mathcal{G} and \mathcal{H} :

$$\widehat{\mathfrak{R}}_{\mathcal{D}}(\mathcal{G}) \leq \frac{1}{2} \widehat{\mathfrak{R}}_{\mathcal{D}_{\mathcal{X}}}(\mathcal{H})$$

• Proposition 2.3 (Rademacher Complexity Generalization Bounds)[Mohri et al., 2018] Let \mathcal{H} be a family of functions taking values in $\{-1, +1\}$ and let \mathcal{P}_X be the distribution over the input space \mathcal{X} . Then, for any $\delta > 0$, with probability at least $1 - \delta$ over a sample \mathcal{D}_m of size m drawn according to \mathcal{P}_X , each of the following holds for any $h \in \mathcal{H}$:

$$L(h) \le \widehat{L}_m(h) + \mathfrak{R}_m(\mathcal{H}) + \sqrt{\frac{\log(1/\delta)}{2m}}$$
 (4)

and

$$L(h) \le \widehat{L}_m(h) + \widehat{\mathfrak{R}}_{\mathcal{D}_m}(\mathcal{H}) + 3\sqrt{\frac{\log(2/\delta)}{2m}}.$$
 (5)

• Remark (Compute Emipirical Radmatcher Complexity) Note that the Radmatcher variable σ_i and $-\sigma_i$ have the same distribution. So

$$\widehat{\mathfrak{R}}_{\mathcal{D}}(\mathcal{H}) = \mathbb{E}_{\sigma} \left[\sup_{h \in \mathcal{H}} \frac{1}{m} \sum_{i=1}^{m} (-\sigma_i h(X_i)) \right] = -\mathbb{E}_{\sigma} \left[\inf_{h \in \mathcal{H}} \frac{1}{m} \sum_{i=1}^{m} \sigma_i h(X_i) \right]$$

Now, for a fixed value of σ , computing

$$\inf_{h \in \mathcal{H}} \frac{1}{m} \sum_{i=1}^{m} \sigma_i h(X_i)$$

is equivalent to an *empirical risk minimization problem*, which is known to be *computationally hard* for some hypothesis sets. Thus, in some cases, computing $\widehat{\mathfrak{R}}_{\mathcal{D}}(\mathcal{H})$ could be computationally hard.

3 VC-Dimension

3.1 Definition of VC-Dimension

• **Remark** Recall the decomposition of generalization error:

$$L(g_n^*) - L^* = \underbrace{\left(L(g_n^*) - \inf_{g \in \mathcal{H}} L(g)\right)}_{estimation\ error} + \underbrace{\left(\inf_{g \in \mathcal{H}} L(g) - L^*\right)}_{approximation\ error}.$$

where the first term is called *estimation error* and the second term is called approximation error. The definition of *PAC learnability* requires that *the estimation error* would be bounded *uniformly over all distributions*.

Can infinite-size classes be learnable, and, if so, what determines their sample complexity?

• Remark (Restriction based on Behavior of Functions on Data D)

Recall the No-Free-Lunch theorem and its proof [Shalev-Shwartz and Ben-David, 2014]. There, we have shown that without restricting the hypothesis class, for any learning algorithm, an adversary can construct a distribution for which the learning algorithm will perform poorly, while there is another learning algorithm that will succeed on the same distribution. To do so, the adversary used a finite set $\mathcal{D} \subset \mathcal{X}$ and considered a family of distributions that are **concentrated on elements of** \mathcal{D} . Each distribution was derived from a "true" target function from \mathcal{D} to $\{0,1\}$. To make any algorithm fail, the adversary used the power of choosing a target function from the set of all possible functions from \mathcal{D} to $\{0,1\}$.

When considering PAC learnability of a hypothesis class \mathcal{H} , the adversary is restricted to constructing distributions for which some hypothesis $h \in \mathcal{H}$ achieves a zero risk. Since we are considering distributions that are concentrated on elements of \mathcal{D} , we should study how \mathcal{H} behaves on \mathcal{D}

• Definition (Restriction of \mathcal{H} to \mathcal{D}).

Let \mathcal{H} be a class of functions from \mathcal{X} to $\{0,1\}$ and let $\mathcal{D} = \{x_1,\ldots,x_m\} \subset \mathcal{X}$.

<u>The restriction of \mathcal{H} to \mathcal{D} is the set of functions from \mathcal{D} to $\{0,1\}$ that can be derived from \mathcal{H} . That is,</u>

$$\mathcal{H}_{\mathcal{D}} := \left\{ \left(h(x_1), \dots, h(x_m) \right) : h \in \mathcal{H} \right\},\,$$

where we **represent** each function from \mathcal{X} to $\{0,1\}$ as a **vector** in $\{0,1\}^{|\mathcal{D}|}$.

• Remark (What You See Is All You Know)

Using the output of functions on a finite set of samples, we can define an equivalence relationship: $f \sim g$ if and only if their outputs vector on given finite set \mathcal{D}_m are the same. Thus, unlike the original space \mathcal{H} , the quotient space \mathcal{H}/\sim is a much simpler function space with finite dimensional representation.

In other word, <u>what you see is all you know</u>, i.e. there is no way to distinguish f and g beyond their answers to given limited set of questions in \mathcal{D} .

• Definition (Shattering).

A hypothesis class \mathcal{H} shatters a finite set $\mathcal{D} \subset \mathcal{X}$ if the restriction of \mathcal{H} to \mathcal{D} is the set

of all functions from \mathcal{D} to $\{0,1\}$. That is,

$$|\mathcal{H}_{\mathcal{D}}| = 2^{|\mathcal{D}|}.$$

- Remark Whenever some set \mathcal{D} is *shattered* by \mathcal{H} , the *adversary* is *not restricted* by \mathcal{H} , as they can *construct a distribution over* \mathcal{D} based on *any* target function from \mathcal{D} to $\{0,1\}$, while still maintaining the realizability assumption.
- The following is the corollary of the No Free Lunch Theorem:

Corollary 3.1 [Shalev-Shwartz and Ben-David, 2014]

Let \mathcal{H} be a hypothesis class of functions from \mathcal{X} to $\{0,1\}$. Let m be a training set size. Assume that there exists a set $\mathcal{D} \subset X$ of size 2m that is **shattered** by \mathcal{H} . Then, for any learning algorithm, \mathcal{A} , there exist a **distribution** \mathcal{P} over $\mathcal{X} \times \{0,1\}$ and a predictor $h \in \mathcal{H}$ such that

$$L_{\mathcal{P}}(h) = 0$$

but with probability of at least 1/7 over the choice of $S \sim P^m$ we have that

$$L_{\mathcal{P}}(\mathcal{A}(\mathcal{S})) \geq 1/8.$$

• Remark (A Model that can Explain Everything is Worthless)
If \mathcal{H} shatters some set \mathcal{D} of size 2m then we cannot learn \mathcal{H} using m examples.

Intuitively, if a set \mathcal{D} is **shattered** by \mathcal{H} , and we receive a sample containing half the instances of \mathcal{D} , the labels of these instances give us **no** information about the labels of the **rest** of the instances in \mathcal{D} - **every possible labeling of the rest of the instances can be explained** by some hypothesis in \mathcal{H} .

Philosophically,

If someone can explain every phenomenon, his explanations are worthless.

• Definition (*VC-Dimension*).

<u>The VC-dimension</u> of a hypothesis class \mathcal{H} , denoted $VCdim(\mathcal{H})$ or simply $v(\mathcal{H})$, is **the** <u>maximal size</u> of a set $\mathcal{D} \subset \mathcal{X}$ that can be **shattered** by \mathcal{H} .

If \mathcal{H} can shatter sets of *arbitrarily large size* we say that \mathcal{H} has *infinite VC-dimension*.

• Theorem 3.2 (No Free Lunch, VC Dimension) [Shalev-Shwartz and Ben-David, 2014] Let \mathcal{H} be a class of infinite VC-dimension. Then, \mathcal{H} is not PAC learnable.

3.2 Growth Function

- Remark We defined the notion of *shattering*, by considering *the restriction of* \mathcal{H} *to a finite set of instances.* The *growth function* measures the *maximal "effective"* size of \mathcal{H} on a set of m examples. Formally:
- Definition (*Growth Function*). Let \mathcal{H} be a hypothesis class. Then <u>the growth function of \mathcal{H} </u>, denoted $\tau_{\mathcal{H}} : \mathbb{N} \to \mathcal{N}$, is defined as

$$\tau_{\mathcal{H}}(m) := \max_{\mathcal{D} \subset \mathcal{X}: |\mathcal{D}| = m} |\mathcal{H}_{\mathcal{D}}|.$$

In words, $\tau_{\mathcal{H}}(m)$ is **the number of different functions** from a set \mathcal{D} of **size** m to $\{0,1\}$ that can be obtained by **restricting** \mathcal{H} **to** \mathcal{D} .

- Remark if $VCdim(\mathcal{H}) = d$ then for any $m \leq d$ we have $\tau_{\mathcal{H}}(m) = 2^m$. In such cases, \mathcal{H} induces all possible functions from \mathcal{D} to $\{0,1\}$.
- Lemma 3.3 (Sauer's Lemma). [Shalev-Shwartz and Ben-David, 2014, Mohri et al., 2018] Let \mathcal{H} be a hypothesis class with $VCdim(\mathcal{H}) \leq d < \infty$. Then, for all m,

$$\tau_{\mathcal{H}}(m) \le \sum_{i=0}^{d} \binom{m}{i} \tag{6}$$

In particular, if m > d + 1 then

$$\tau_{\mathcal{H}}(m) \le \left(\frac{e \, m}{d}\right)^d.$$
(7)

Proof: To prove the lemma it suffices to prove the following *stronger claim*: For any $\mathcal{D} = \{x_1, \ldots, x_m\}$ we have

$$\forall \mathcal{H}, |\mathcal{H}_{\mathcal{D}}| \le |\{\mathcal{B} \subseteq \mathcal{D} : \mathcal{H} \text{ shatters } \mathcal{B}\}|$$
 (8)

The reason why Equation (8) is sufficient to prove the lemma is that if $VCdim(\mathcal{H}) \leq d$ then no set whose size is larger than d is shattered by \mathcal{H} and therefore

$$|\{\mathcal{B} \subseteq \mathcal{D} : \mathcal{H} \text{ shatters } \mathcal{B}\}| \leq \sum_{i=0}^{d} {m \choose i}.$$

When m > d + 1 the right-hand side of the preceding is at most $\left(\frac{e m}{d}\right)^d$.

We are left with proving Equation (8) and we do it using an *inductive argument*.

- 1. For m = 1, no matter what \mathcal{H} is, either both sides of Equation (8) equal 1 or both sides equal 2 (the empty set is always considered to be shattered by \mathcal{H}).
- 2. Assume Equation (8) holds for sets of size k < m and let us prove it for sets of size m.

Fix \mathcal{H} and $\mathcal{D} = \{x_1, \dots, x_m\}$. Denote $\mathcal{D}_{-1} = \{x_2, \dots, x_m\}$ and in addition, define the following two sets:

$$Y_0 = \{(y_2, \dots, y_m) : (0, y_2, \dots, y_m) \in \mathcal{H}_{\mathcal{D}} \lor (1, y_2, \dots, y_m) \in \mathcal{H}_{\mathcal{D}} \},$$

and

$$Y_1 = \{(y_2, \dots, y_m) : (0, y_2, \dots, y_m) \in \mathcal{H}_{\mathcal{D}} \land (1, y_2, \dots, y_m) \in \mathcal{H}_{\mathcal{D}}\}.$$

It is easy to verify that $|\mathcal{H}_{\mathcal{D}}| = |Y_0| + |Y_1|$. Additionally, since $Y_0 = \mathcal{H}_{\mathcal{D}_{-1}}$, using the induction assumption (applied on \mathcal{H} and \mathcal{D}_{-1}) we have that

$$|Y_0| = |\mathcal{H}_{\mathcal{D}_{-1}}| \le |\{\mathcal{B} \subseteq \mathcal{D}_{-1} : \mathcal{H} \text{ shatters } \mathcal{B}\}|$$

= $|\{\mathcal{B} \subseteq \mathcal{D} : x_1 \notin \mathcal{B} \land \mathcal{H} \text{ shatters } \mathcal{B}\}|.$

Next, define $\mathcal{H}' \subseteq \mathcal{H}$ to be

$$\mathcal{H}' = \left\{ h \in \mathcal{H} : \exists h' \in \mathcal{H} \text{ s.t. } (1 - h'(x_1), h'(x_2), \dots, h'(x_m)) = (h(x_1), h(x_2), \dots, h(x_m)) \right\},$$

namely, \mathcal{H}' contains pairs of hypotheses (h, h') that agree on \mathcal{D}_{-1} and differ on x_1 . Using this definition, it is clear that if \mathcal{H}' shatters a set $\mathcal{B} \subseteq \mathcal{D}$ then it also shatters the set $\mathcal{B} \cup \{x_1\}$ and vice versa.

Combining this with the fact that $Y_1 = \mathcal{H}'_{\mathcal{D}_{-1}}$ and using the inductive assumption (now applied on \mathcal{H}' and \mathcal{D}_{-1}) we obtain that

$$|Y_{1}| = |\mathcal{H}'_{\mathcal{D}_{-1}}| \le |\{\mathcal{B} \subseteq \mathcal{D}_{-1} : \mathcal{H}' \text{ shatters } \mathcal{B}\}|$$

$$= |\{\mathcal{B} \subseteq \mathcal{D}_{-1} : \mathcal{H}' \text{ shatters } \mathcal{B} \cup \{x_{1}\}\}|$$

$$= |\{\mathcal{B} \subseteq \mathcal{D} : x_{1} \in \mathcal{B} \wedge \mathcal{H}' \text{ shatters } \mathcal{B}\}|$$

$$= |\{\mathcal{B} \subseteq \mathcal{D} : x_{1} \in \mathcal{B} \wedge \mathcal{H} \text{ shatters } \mathcal{B}\}|.$$

Overall, we have shown that

$$\begin{aligned} |\mathcal{H}_{\mathcal{D}}| &= |Y_0| + |Y_1| \\ &\leq |\{\mathcal{B} \subseteq \mathcal{D} : x_1 \notin \mathcal{B} \land \mathcal{H} \text{ shatters } \mathcal{B}\}| + |\{\mathcal{B} \subseteq \mathcal{D} : x_1 \in \mathcal{B} \land \mathcal{H} \text{ shatters } \mathcal{B}\}| \\ &= |\{\mathcal{B} \subseteq \mathcal{D} : \mathcal{H} \text{ shatters } \mathcal{B}\}| \end{aligned}$$

which concludes our proof.

3.3 Relate Growth Function to Radmatcher Complexity

• Lemma 3.4 (Massart's Lemma) [Mohri et al., 2018] Let $A \subseteq \mathbb{R}^m$ be a finite set, with $r = \max_{x \in A} ||x||_2$, then the following holds:

$$\mathbb{E}_{\sigma} \left[\frac{1}{m} \sup_{x \in A} \sum_{i=1}^{m} \sigma_i x_i \right] \le \frac{r\sqrt{2\log|A|}}{m} \tag{9}$$

where σ_i 's are independent uniform random variables taking values in $\{-1, +1\}$ and x_1, \ldots, x_m are the components of vector x.

• Corollary 3.5 (Radmatcher Complexity Bounds by Growth Number) [Mohri et al., 2018]

Let \mathcal{H} be a family of functions taking values in $\{-1, +1\}$. Then the following holds:

$$\mathfrak{R}_m(\mathcal{H}) \le \sqrt{\frac{2\log \tau_{\mathcal{H}}(m)}{m}} \tag{10}$$

Proof: For a fixed sample $\mathcal{D} = (x_1, \dots, x_m)$, we denote by $\mathcal{H}_{\mathcal{D}} \subset \{-1, +1\}^m$ the set of vectors of function values $(h(x_1), \dots, h(x_m))$ where h is in \mathcal{H} . Since $h \in \mathcal{H}$ takes values in $\{-1, +1\}$, the *norm* of these vectors is bounded by \sqrt{m} . We can then apply *Massart's lemma* as follows:

$$\mathfrak{R}_{m}(\mathcal{H}) = \mathbb{E}_{\mathcal{D}} \left[\mathbb{E}_{\sigma} \left[\sup_{u \in \mathcal{H}_{\mathcal{D}}} \frac{1}{m} \sum_{i=1}^{m} \sigma_{i} u_{i} \middle| \mathcal{D} \right] \right] \leq \mathbb{E}_{\mathcal{D}} \left[\frac{\sqrt{m} \sqrt{2 \log |\mathcal{H}_{\mathcal{D}}|}}{m} \right]$$

By definition, $|\mathcal{H}_{\mathcal{D}}|$ is bounded by the growth function, thus,

$$\mathfrak{R}_m(\mathcal{H}) \leq \mathbb{E}_{\mathcal{D}}\left[\frac{\sqrt{m}\sqrt{2\log\tau_{\mathcal{H}}(m)}}{m}\right] = \sqrt{\frac{2\log\tau_{\mathcal{H}}(m)}{m}},$$

which concludes the proof.

3.4 Generalization Bounds via Growth Function and VC-Dimension

• Combining Proposition 2.3 to Corollary 3.5, we have:

Corollary 3.6 (Growth Function Generalization Bound) [Mohri et al., 2018] Let \mathcal{H} be a family of functions taking values in $\{-1, +1\}$. Then, for any $\delta > 0$, with probability at least $1 - \delta$, for any $h \in \mathcal{H}$,

$$L(h) \le \widehat{L}_m(h) + \sqrt{\frac{2\log \tau_{\mathcal{H}}(m)}{m}} + \sqrt{\frac{\log(1/\delta)}{2m}}$$
(11)

Growth function bounds can be also derived directly (without using Rademacher complexity bounds first). The resulting bound is then the following:

$$\mathcal{P}\left\{ \left| L(h) - \widehat{L}_m(h) \right| > \epsilon \right\} \le 4\tau_{\mathcal{H}}(2m) \exp\left(-\frac{m\epsilon^2}{8}\right)$$
 (12)

which only differs from (11) by constants.

• Applying Sauer's Lemma to Corollary 3.6, we have:

Corollary 3.7 (VC-Dimension Generalization Bounds) [Mohri et al., 2018] Let \mathcal{H} be a family of functions taking values in $\{-1, +1\}$ with VC-dimension d. Then, for any $\delta > 0$, with probability at least $1 - \delta$, the following holds for all $h \in \mathcal{H}$:

$$L(h) \le \widehat{L}_m(h) + \sqrt{\frac{2d\log(em/d)}{m}} + \sqrt{\frac{\log(1/\delta)}{2m}}$$
(13)

Thus, the form of this generalization bound is

$$L(h) \le \widehat{L}_m(h) + O\left(\sqrt{\frac{\log(m/d)}{(m/d)}}\right),$$
 (14)

which emphasizes the importance of the ratio m/d for generalization.

• Remark The theorem provides another instance of <u>Occam's razor principle</u> where simplicity is <u>measured</u> in terms of <u>smaller VC-dimension</u>.

VC-dimension bounds can be derived directly without using an intermediate Rademacher complexity bound: combining Sauers lemma with (12) leads to the following high-probability bound

$$L(h) \le \widehat{L}_m(h) + \sqrt{\frac{8d\log(2em/d) + 8\log(4/\delta)}{m}},\tag{15}$$

which has the general form of (14). The log factor plays only a minor role in these bounds. A finer analysis can be used in fact to eliminate that factor.

3.5 Examples

3.6 Lower Bounds

• Remark (No-Free-Lunch via VC-Dimension)
This section provides lower bounds on the generalization error of any learning algorithm in terms of the VC-dimension of the hypothesis set used.

These lower bounds are shown by finding for any algorithm a 'bad' distribution. Since the learning algorithm is arbitrary, it will be difficult to specify that particular distribution. Instead, it suffices to prove its existence non-constructively. At a high level, the proof technique used to achieve this is the probabilistic method of Paul Erdös.

In the context of the following proofs,

- 1. first a lower bound is given on the expected error over the parameters defining the distributions.
- 2. From that, the lower bound is shown to hold for at least one set of parameters, that is one distribution.
- Proposition 3.8 (Lower Bound, Deterministic Case) [Mohri et al., 2018] Let \mathcal{H} be a hypothesis set with VC-dimension d > 1. Then, for any learning algorithm \mathcal{A} , there exist a distribution \mathcal{P} over \mathcal{X} and a target function $f \in \mathcal{H}$ such that

$$\mathcal{P}_{\mathcal{D}_m}\left\{L_{\mathcal{P}}(g_m, f) > \frac{d-1}{32m}\right\} \ge \frac{1}{100}$$

• Remark (Infinite VC-dimension \Rightarrow Non PAC-learnable)

The theorem shows that for any algorithm \mathcal{A} , there exists a 'bad' distribution over \mathcal{X} and a target function f for which the error of the hypothesis returned by \mathcal{A} is

$$\Omega\left(\frac{d}{m}\right)$$

with some constant probability. This further demonstrates the key role played by the VC-dimension in learning. The result implies in particular that PAC-learning in the non-realizable case is not possible when the VC-dimension is infinite.

• Proposition 3.9 (Lower Bound, Non-Realizable (Stochastic) Case) [Mohri et al., 2018]

Let \mathcal{H} be a hypothesis set with **VC-dimension** d > 1. Then, for any learning algorithm \mathcal{A} , there exists a distribution \mathcal{P} over $\mathcal{X} \times \{0,1\}$ such that:

$$\mathcal{P}_{\mathcal{D}_m}\left\{L_{\mathcal{P}}(g_m(\cdot|\mathcal{D}_m)) - \inf_{g \in \mathcal{H}} L_{\mathcal{P}}(g) > \sqrt{\frac{d}{320m}}\right\} \ge \frac{1}{64}$$

Equivalently, for any learning algorithm, the sample complexity verifies

$$m \ge \frac{d}{320\epsilon^2}.$$

• Remark The theorem shows that for any algorithm \mathcal{A} , in the non-realizable case (i.e. the Bayes classifier is not in \mathcal{H}), there exists a 'bad' distribution over $\mathcal{X} \times \{0,1\}$ such that the error of the hypothesis returned by \mathcal{A} is

$$\Omega\left(\sqrt{\frac{d}{m}}\right)$$

with some constant probability. The *VC-dimension* appears as a critical quantity in learning in this general setting as well. In particular, with an *infinite VC-dimension*, agnostic *PAC-learning is not possible*.

References

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