

# Philosophy 12, Problem Set 6

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## 1 Q1

### 1.1 a

Randomly pick items into the bag until we overshoot the maximum weight  $W$ . Then remove the last item and check if what's in the bag has at least  $V$  calories. If it does, we have shown that  $S$  exists. If not we take every item out and repeat the algorithm.

This is not a polynomial time algorithm. Since we are randomly selecting the items, it is possible to never find such a set  $S$  even if  $S$  exists.

### 1.2 b

Since each item can either be picked or not picked, there are  $2^n$  ways to pick our items for the backpack. Now we can simply iterate through all  $2^n$  options to see if any fits the requirement of having less weight than  $W$  and more calories than  $V$ .

This is also not a polynomial time algorithm since  $2^n$  is exponential.

## 2 Q2

### 2.1 Each country has at least one color

Let  $i$  be an arbitrary country, then we have  $c_i \vee m_i \vee y_i$ . Therefore we can apply to all  $i$ :

$$\bigwedge_{i=1}^n (c_i \vee m_i \vee y_i)$$

## 2.2 Each country has at most one color

Let  $i$  be an arbitrary country, then we have  $(c_i \wedge \neg m_i \wedge \neg y_i) \vee (m_i \wedge \neg y_i \wedge \neg c_i) \vee (y_i \wedge \neg c_i \wedge \neg m_i)$ . Therefore we can apply to all  $i$ :

$$\bigwedge_{i=1}^n (c_i \wedge \neg m_i \wedge \neg y_i) \vee (m_i \wedge \neg y_i \wedge \neg c_i) \vee (y_i \wedge \neg c_i \wedge \neg m_i)$$

## 2.3 No adjacent countries have the same color

If two adjacent countries have the same color then  $(c_i \wedge c_j) \vee (m_i \wedge m_j) \vee (y_i \wedge y_j)$   
We can negate that for all such  $i, j$

$$\begin{aligned} & \neg \bigvee_{i,j \text{ adjacent}} (c_i \wedge c_j) \vee (m_i \wedge m_j) \vee (y_i \wedge y_j) \\ & \equiv \bigwedge_{i,j \text{ adjacent}} \neg(c_i \wedge c_j) \wedge \neg(m_i \wedge m_j) \wedge \neg(y_i \wedge y_j) \end{aligned}$$

Finally we can simply combine these, so our answer is  $(2.1) \wedge (2.2) \wedge (2.3)$

## 3 Q5

The first statement is true because there are only two Republicans in the race: Reagan and Anderson. The second statement is where the problem begins. “A Republican” will win the election is true in the pollee’s mind because Regan is a Republican. They have subsituted “a republican” with “Reagan” because he is decisively ahead of Carter. Therefore the final conclusion: “If it’s not Regan who wins, it will be Anderson” doesn’t make sense.

This shows that everyday English is not as strict as formal logic and contains a lot of implicit substitutions of terms.

1 | Lemma 1

2 | Lemma 2

3 | Lemma 3

4 |  $C$  is real

5 |  $C-1$  is real  $(E, 1, 4)$

6 |  $0 \leq (C-1)^2$   $(E, 2, 5)$

7. |  $0 \leq C^2 + 1 - 2C$

8. |  $2C \leq C^2 + 1$

9. |  $C$  is real  $\rightarrow 2C \leq C^2 + 1$   $(I, 4, 8)$

(a)

1.		<u>P</u>	
2.			<u>q</u>
3.			q (R 2)
4.			$q \rightarrow q$ (I, 2, 3)
5.			$P \rightarrow (q \rightarrow q)$ (I, 1, 4)

(b)  $r \rightarrow (p \rightarrow (q \rightarrow p))$

1.		<u>r</u>			
2.			<u>p</u>		
3.				<u>q</u>	
4.				p	(R, 2)
5.				$q \rightarrow p$	(I, 3, 4)
6.				$p \rightarrow (q \rightarrow p)$	(I, 2, 5)
7.				$r \rightarrow (p \rightarrow (q \rightarrow p))$	(I, 1, 6)

(c)  $(p \rightarrow q) \rightarrow ((p \rightarrow (q \rightarrow r)) \rightarrow (p \rightarrow r))$

1. | |  $p \rightarrow q$

1.	$P \rightarrow q$	
2.	$P \rightarrow (q \rightarrow r)$	
3.	$P$	
4.	$q$	(E, 1, 3)
5.	$q \rightarrow r$	(E, 2, 3)
6.	$P \rightarrow r$	(I, 3, 5)
7.	$(P \rightarrow (q \rightarrow r)) \rightarrow (P \rightarrow r)$	(I, 2, 6)
8.	$(p \rightarrow q) \rightarrow ((P \rightarrow (q \rightarrow r)) \rightarrow (P \rightarrow r))$	(I, 1, 7)

(d)  $(p \rightarrow (q \rightarrow r)) \rightarrow (q \rightarrow (p \rightarrow r))$

1.	$P \rightarrow (q \rightarrow r)$	
2.	$q$	
3.	$P$	
4.	$q \rightarrow r$	(E, 1, 3)
5.	$r$	(E, 2, 4)
6.	$P \rightarrow r$	(I, 3, 5)
7.	$q \rightarrow (P \rightarrow r)$	(I, 2, 6)
8.	$(P \rightarrow (q \rightarrow r)) \rightarrow (q \rightarrow (P \rightarrow r))$	(I, 1, 7)

$$8. \mid (p \rightarrow (q \rightarrow r)) \rightarrow (q \rightarrow (p \rightarrow r)) \quad (I, 1, 7)$$

$$(e) \quad (p \rightarrow (q \rightarrow r)) \rightarrow ((p \rightarrow q) \rightarrow (p \rightarrow r))$$

1.	$p \rightarrow (q \rightarrow r)$	
2.	$(p \rightarrow q)$	
3.	$p$	
4.	$q$	$(E, 2, 3)$
5.	$q \rightarrow r$	$(E, 1, 3)$
6.	$r$	$(E, 4, 5)$
7.	$p \rightarrow r$	$(I, 3, 6)$

$$8. \mid (p \rightarrow q) \rightarrow (p \rightarrow r) \quad (I, 2, 7)$$

$$9. \mid (p \rightarrow (q \rightarrow r)) \rightarrow ((p \rightarrow q) \rightarrow (p \rightarrow r)) \quad (I, 1, 8)$$