

# Lab 1: My Malloc Library

CS-350: Systems Programming

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# my\_malloc library

`void * my_malloc(uint32_t size)`  
allocates size bytes of memory

`void my_free(void *ptr)`  
deallocates memory pointed to by `ptr`, previously allocated by `my_malloc()`

`void coalesce_free_list(void)`  
merges adjacent chunks on the free list into single larger chunks

`FreeListNode * free_list_begin( void )`  
retrieves the first node of the free list

# Memory Allocation

```
void * my_malloc(uint32_t sz):
```

1. Find chunk (from heap or free list)
2. Split chunk if too large
  - put remainder on free list
3. Chunk to return
  1. Header (for bookkeeping)
    - 1<sup>st</sup> 4-bytes: total chunk size (including header+fragment)
    - 2<sup>nd</sup> 4-bytes: magic number (validates malloc'd chunks)
  2. User allocation: `sz` bytes
    - Return chunk ptr: 8-bytes inside chunk (after header)
  3. Fragmentation:
    - Any necessary padding or wastage from oversized chunk

0xffff000  
(header: 8 bytes)

4-byte chunk size  
4-byte checksum

0xffff008  
(return ptr)

User allocation  
(`sz` bytes)

0xffff008+sz

Fragmentation  
(padding + oversized wastage)

# Free List Management

`my_free( ptr )` places chunk on free list

```
typedef struct freelistnode {  
    struct freelistnode *flink;  
    uint32_t size;  
} * FreeListNode;
```

1. Check for valid magic number
2. Embed `struct freelistnode` at start of chunk
3. Insert free list node into free list



# Auxiliary Functions

```
void coalesce_free_list(void)
```

Merges adjacent chunks on the free list into single larger chunks

Free list nodes should be kept in order

No coalescing unless this function is called!

```
FreeListNode * free_list_begin( void )
```

retrieves the first node of the free list

# Requirements and Constraints

1. **sbrk()** is the only allowed third party library or system call allowed
2. Always call `sbrk(8192)` except if `my_malloc()` needs more bytes
3. Assume that other library routines also may call `sbrk()`.
4. You may not use more than 8 bookkeeping bytes.
5. You may use **one** global variable for the first free list node
6. Free list should always be sorted in ascending order by chunk address.
7. Use *first fit* strategy to search the free list, i.e. return the first usable chunk.

You may call `sbrk(0)` to identify the heap's current end.

# What I did! (Not necessarily what you have to do)

## 1. Implemented malloc()

1. find\_chunk(): returns address of appropriately sized chunk to use
  - first only from the heap, then later from the free list first then heap
2. split\_chunk(): returns address of chunk to use
  - if needed, split chunk and put remainder on free list
3. Bookkeep: place chunk size and magic number in header
4. Return ptr to user allocation

## 2. Implemented free()

1. implemented singly-linked list for `struct freelistnode`
2. check for magic number
3. embed struct freelistnode into chunk
4. insert chunk into free list

## 3. Implemented coalesce()

1. wrote function to test if two nodes are adjacent
2. wrote function to merge two adjacent nodes
3. traverse free list testing and merging adjacent nodes

# Other Hints and Tips

- Start early!
- Don't start programming until you fully understand the concepts:
  - this lab is complex, but not a lot of code
- Remember, pointer arithmetic is based on pointer type
- Build and test incrementally
- Consider a driver program that validates heap and free list after each malloc()/free()
- Don't forget "my\_errno"
- Memory debuggers (gdb, valgrind, etc.) are your friends