MAT 344 Lecture Notes

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`Github Page https://github.com/TianyuDu/Spikey_UofT_Notes`
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1 Strings, Sets, and Binomial Coefficients

1.1 Strings and Sets

Notation 1.1. Let $n \in \mathbb{Z}_{++}$, and we use $[n]$ to denote the n-element set $\{1, 2, \ldots, n\}$.

Definition 1.1. Let X be a set, then an X-**string of length** (or a **word/array**) n is a function $s : [n] \to X$, and X is called the **alphabet** of the string, and each $x \in X$ is called a **character** or letter.

Remark 1.1. An X-string defined by $s : [n] \to X$ with length n can be equivalently defined as a **sequence** consisting elements in X.

$$s(1)s(2)\ldots s(n) \tag{1.1}$$

Definition 1.2. In the case $X = \{0, 1\}$, strings generated from X are called **binary strings**. When $X = \{0, 1, 2\}$, strings are called **ternary strings**.

Definition 1.3. Let X be a *finite* set and let $n \in \mathbb{Z}_{++}$. An X-string $s = x_1x_2\ldots x_n$ is a **permutation** of size m if $x_i \neq x_j\ \forall x_i, x_j \in s$.

Proposition 1.1. If X is an m-element set and $m \geq n \in \mathbb{Z}_{++}$, then the number of X-strings of length n that are permutations is

$$P(m, n) \equiv \frac{m!}{(m-n)!} \tag{1.2}$$

Definition 1.4. Let X be a *finite* set and let $0 \leq k \leq |X|$. Then $S \subseteq X$ with $|S| = k$ is a **combination** of size k.

Proposition 1.2. Let $n, k \in \mathbb{Z}$ such that $0 \leq k \leq n$, then the number of combinations is

$$\binom{n}{k} \equiv \frac{P(n,k)}{n!} = \frac{n!}{k!(n-k)!} \tag{1.3}$$

Proposition 1.3. For all integers n and k with $0 \leq k \leq n$

$$\binom{n}{k} = \binom{n}{n-k} \tag{1.4}$$

Example 1.1. Binomial coefficients can be used to find the number of integer solutions of

$$\sum_{i=1}^{k} x_i \leq N \tag{1.5}$$

given appropriate integers $k, N \in \mathbb{Z}$.

(i) $x_i > 0\ \forall i \in [k]$ and equality holds, then $C(N-1, k-1)$.

(ii) $x_i \geq 0\ \forall i \in [k]$ and equality holds, then $C(N+k-1, k-1)$.[1]

(iii) $x_i > 0\ \forall i \neq j, x_j = Z$ and equality holds, then $C(N-Z+k-2, k-2)$.

(iv) $x_i > 0\ \forall i \in [k]$ and strict inequality holds, then $C(N-1, k)$.[2]

(v) $x_i \geq 0\ \forall i \in [k]$ and strict inequality holds, then $C(N+k-1, k)$.

(vi) $x_i \geq 0\ \forall i \in [k]$ and *weak* inequality holds, $C(N+k, k)$ [3].

$$\binom{N+k-1}{k-1} + \binom{N+k-1}{k} = \binom{N+k}{k} \tag{1.6}$$

[1] Simulate choosing $x_i + 1$ instead of x_i.
[2] Image there is a placeholder $x_{k+1} > 0$.
[3] This can be calculated by adding case (ii) and case (v) together, and apply Pascal's identity

Definition 1.5. Define a **plane** as $\mathbb{Z}^2$, then a **lattice path** in the plane is a *sequence* of elements in $\mathbb{Z}^2$

$$((x_i, y_i))_{i=1}^t \tag{1.7}$$

such that for every $i \in \{1, \ldots, t-1\}$, either

(i) (*Horizontal move*) $x_{i+1} = x_i + 1 \wedge y_{i+1} = y_i$

(ii) Or (*vertical move*) $x_{i+1} = x_i \wedge y_{i+1} = y_i + 1$

Lemma 1.1. Let $(p, q), (m, n) \in \mathbb{Z}^2$, then the number of lattice paths from (p, q) to (m, n) is

$$\binom{(p-m) + (q-n)}{p-m} \tag{1.8}$$

Proof. The lattice is isomorphic to a H, V-string with length $(p-m) + (q-n)$. There are exactly $p - m$ horizontal moves as well as exactly $q - n$ vertical moves. ■

Theorem 1.1. Given $n \in \mathbb{Z}_+$, the number of lattice paths from $(0, 0)$ to (n, n) which *never go above the diagonal line* is the **Catalan number**

$$C(n) \equiv \frac{1}{n+1}\binom{2n}{n} \tag{1.9}$$

Proof. Omitted ■

Theorem 1.2 (Binomial Theorem)**.** Let $x, y \in \mathbb{R}$, then $\forall n \in \mathbb{Z}_+$

$$(x+y)^n = \sum_{i=0}^{n} \binom{n}{i} x^{n-i} y^i \tag{1.10}$$

Theorem 1.3 (Multinomial Theorem)**.** Let $r \in \mathbb{Z}_+$, $\{x_i\}_{i=1}^r \in \mathcal{P}(\mathbb{R})$. Then for every $n \in \mathbb{Z}_+$,

$$(\sum_{i=1}^{r} x_i)^n = \sum_{|\alpha|=n} \binom{n}{\alpha} (x_i)^\alpha \tag{1.11}$$

where $\alpha \equiv (\alpha_i)_{i=1}^r,\ \alpha_i \in \mathbb{Z}_{++}\ \forall i$ is a **multi-index**, and

$$(x_i)^\alpha \equiv \sum_{i=1}^{r} x_i^{\alpha_i} \tag{1.12}$$

$$|\alpha| \equiv \sum_{i=1}^{r} \alpha_i \tag{1.13}$$

$$\binom{n}{\alpha} \equiv \frac{n!}{\alpha_1!\alpha_2!\ldots\alpha_r!} \tag{1.14}$$

2 Induction

Theorem 2.1 (Well-Ordering Principle)**.** Every non-empty set of Z_{++} has a least element.

Proof. Prove using principle of mathematical induction and contradiction. ■

Definition 2.1. Recursive definition

Theorem 2.2 (The Principle of Mathematical Induction)**.** If S is any set of natural numbers with properties that

1. 1 is in S, and

2. $k+1$ is in S whenever k is any number in S.

then $S = \mathbb{Z}_+$.

Remark 2.1. **Recursive definitions** can also be recast as **inductive definitions**.

Definition 2.2 (Summation)**.** Summation operator beginning with index 1, $\sum : \mathcal{F}_1 \times Z_{++} \to \mathbb{R}$, where $\mathcal{F}_1$ is the set of unary real-valued functions, is defined inductively as

$$\sum_{i=1}^{1} f(i) \equiv f(1) \tag{2.1}$$

$$\sum_{i=1}^{k+1} f(i) \equiv \sum_{i=1}^{k} f(i) + f(k+1) \tag{2.2}$$

Theorem 2.3 (The Principle of Complete Mathematical Induction)**.** If S is any set of natural numbers with the properties that

1. $1 \in S$, and
2. $\{1, 2, \ldots, k\} \subset S \implies k+1 \in S$,

then $S = \mathbb{Z}_+$.

3 Pigeon Hole Principle and Complexity

3.1 Pigeon Hole Principle

Theorem 3.1. Let $f : X \to Y$ be a function, then

$$f \text{ injective} \implies |X| \leq |Y| \tag{3.1}$$

Theorem 3.2 (Pigeon Hole Principle)**.** Let $f : X \to Y$, and suppose $|X| > |Y|$, then f is not injective, that's

$$\exists x_1 \neq x_2 \in X \ s.t. \ f(x_1) = f(x_2) \tag{3.2}$$

Proof. Contrapositive form of the theorem 3.1 ■

Theorem 3.3 (Erods/Szekeres)**.** Let $m, n \in \mathbb{Z}_+$, then any sequence of $mn+1$ *distinct* real numbers either

(i) has an increasing subsequence of $m+1$ terms,

(ii) or it has a decreasing subsequence of $n+1$ terms.

Proof. Let $\sigma = (x_1, x_2, \ldots, x_{mn+1})$ be a sequence with length $mn+1$ consisting of distinct reals.
For each $i \in [mn+1]$ define a_i as the maximum length of an increasing subsequence of σ *beginning with* x_i.
Define b_i as the maximum length of a decreasing subsequence of σ *ending with* x_i.
Case (i)

$$\exists i \in [mn+1] \ s.t. \ a_i \geq m+1 \vee b_i \geq n+1 \tag{3.3}$$

then the theorem is proven.
Case (ii) Suppose otherwise

$$\forall i \in [mn+1] \ a_i \leq m \wedge b_i \leq n \tag{3.4}$$

construct function $f : [mn+1] \to [m] \times [n]$ defined as

$$f(i) \equiv (a_i, b_i) \tag{3.5}$$

Note that $|[mn+1]| > |[m] \times [n]|$ so f cannot be injective, so there exists $j \neq k \in [mn+1]$ such that $(a_j, b_j) = (a_k, b_k)$.
WLOG, assume $j < k$.
Since all elements in σ are distinct, $j \neq k \implies x_j \neq x_k$.
Sub-case (i) $x_j < x_k$, then any increasing subsequence beginning with x_k can be extended by prepending x_j, so $a_j > a_k$.
Sub-case (ii) $x_j > x_k$, then any decreasing subsequence ending with x_j can be extended by appending x_k, so $b_k > b_j$.
Either sub-case leads to a contradiction, so **case (ii)** is impossible. ■

3.2 Complexity

Definition 3.1. Let $f, g : \mathbb{N} \to \mathbb{R}$ be a function, then the **big oh** $\mathcal{O}(f)$ is a collection of functions such that, for every $g \in \mathcal{O}(f)$

$$\exists c \in \mathbb{R}, n^* \in \mathbb{N} \; s.t. \; \forall n \in \mathbb{N}, n \geq n^* \implies g(n) \leq cf(n) \tag{3.6}$$

Definition 3.2. Let $f, g : \mathbb{N} \to \mathbb{R}$ be a function. If $f(n) > 0 \; \forall n \in \mathbb{N}$, then the **little oh** $o(f)$ is the collection of functions such that, for every $g \in o(f)$,

$$\lim_{n\to\infty} \frac{g(n)}{f(n)} = 0 \tag{3.7}$$

Definition 3.3. Let $f, g : \mathbb{N} \to \mathbb{R}$, then the **little oh**, $o(f)$ is defined as the collection of functions such that $g \in o(f)$ if and only if

$$\exists c \in \mathbb{R}, n^* \in \mathbb{N}, \; s.t. \; \forall n \in \mathbb{N}, n \geq n^* \implies |g(n)| < c|f(n)| \tag{3.8}$$

Definition 3.4. Define $\pi : \mathbb{Z}_{++} \to \mathbb{Z}_{+}$ as $\pi(n) \equiv$ *the number of primes among the first n positive integers.*

Theorem 3.4 (Prime Number Theorem)**.** $\pi(n)$ grows at a rate the same as $\frac{n}{\ln(n)}$. That's

$$\lim_{n\to\infty} \pi(n) \frac{\ln(n)}{n} = 1 \tag{3.9}$$

Definition 3.5. The class of **polynomial time** problems, denoted as $\mathcal{P}$, is the set of decision problems for which there exists one polynomial run time algorithm as the solution.

Definition 3.6. The class of **nondeterministic polynomial time** problems, denoted as $\mathcal{NP}$, is the set of decision problems for which there is a certificate for a yes answer whose correctness can be verified in polynomial time.

4 Graph Theory

Definition 4.1. A graph $\mathcal{G}$ is defined as an order pair of sets (V, E). **Vertex** set V is a set consisting of **vertex** objects. **Edge set** E contains **edges** as pairs of elements in E.

Definition 4.2. A graph $\mathcal{G}$ is called a **simple graph** if it is unweighted, undirected and contains no loop or multiple edges. That's, if $\mathcal{G} \equiv (V, E)$ is a simple graph, then

1. (Undirected) $\forall x, y \in V, \; xy \in E \iff yx \in E$.
2. (No loop) $\forall xy \in E, x \neq y$.
3. (No multiple edge) all elements in E are distinct.
4. Vertices or edges in $\mathcal{G}$ have no weight.

Graphs with multiple edges or loops are called **multi-graphs**.

Remark 4.1. In this course, unless explicitly mentioned, we consider simple graphs only.

Definition 4.3. Let $x, y \in V$, if $xy \in E$, then x and y are **adjacent**, and edge xy is **incident to** vertices x and y. If $xy \notin E$, we say x and y are **non-adjacent**.

Definition 4.4. Let $\mathcal{G} \equiv (V, E)$ and $x \in V$, then the **neighbourhood** of x is defined as

$$\mathcal{N}(x) \equiv \{v \in V : xy \in E\} \tag{4.1}$$

Then the **degree** of x in graph $\mathcal{G}$ is defined as

$$\deg_{\mathcal{G}}(x) \equiv |\mathcal{N}(x)| \tag{4.2}$$

Definition 4.5. Let $\mathcal{G} \equiv (V, E)$ and $\mathcal{H} \equiv (W, F)$, we say $\mathcal{H}$ is a **subgraph** of $\mathcal{G}$ when $W \subseteq V$ and $F \subseteq E$. $\mathcal{H}$ is an **induced subgraph** if

$$F = \{xy \in E : x, y \in W\} \tag{4.3}$$

$\mathcal{H}$ is a **spanning subgraph** if $W = V$.

Definition 4.6. $\mathcal{G} \equiv (V, E)$ is a **complete graph** ($\mathbf{K}_n$) if

$$E = \{xy : \text{ distinct pair } x, y \in V\} \tag{4.4}$$

Definition 4.7. A graph $\mathcal{G} \equiv (V, E)$ is a **independent graph** ($\mathbf{I}_n$) if for every distinct pair $(x, y) \subset V$, $xy \notin E$.

Definition 4.8. A **walk** in graph $\mathcal{G} \equiv (V, E)$ is a *sequence of vertices* $(x_1, x_2, \ldots, x_n)$ such that

$$x_i x_{i+1} \in E \ \forall i \in \{1, \ldots, n-1\} \tag{4.5}$$

Definition 4.9. A **path** is a walk with *distinct* vertices. The length of path is defined as the number of edges in it.

Definition 4.10. A **cycle** is a *path* $(x_1, x_2, \ldots, x_n)$ with $n \neq 3$ such that $x_1 x_n \in E$.

Definition 4.11. Two graphs $\mathcal{G} \equiv (V, E)$ and $\mathcal{H} \equiv (W, F)$ are **isomorphic**, denoted as $\mathcal{G} \cong \mathcal{H}$, if there exists a bijection $f : V \to W$ such that

$$\forall x, y \in V, \ xy \in E \iff f(x)f(y) \in F \tag{4.6}$$

And we say $\mathcal{G}$ *contains* $\mathcal{H}$ is there is a subgraph of $\mathcal{G}$ isomorphic to $\mathcal{H}$.

Definition 4.12. A graph $\mathcal{G}$ is **connected** when for every distinct pair $x, y \in V$, there exists a path from x to y. Otherwise, $\mathcal{G}$ is **disconnected**.

Definition 4.13. Provided $\mathcal{G}$ is disconnected, then a **component** of $\mathcal{G}$ is a *maximal connect subgraph* of $\mathcal{G}$. That's, if $\mathcal{H}$ is a component, it is connected and any super-graph of $\mathcal{H}$ is disconnected.

Definition 4.14. A graph $\mathcal{G}$ is **acyclic** when it does not contain any cycle on three or more vertices. An acyclic graph is also called **forests**. Further, if a acyclic graph is connected, it's called a **tree**.

Definition 4.15. Given $\mathcal{G}$ is connected[4], a subgraph $\mathcal{H} \equiv (W, F)$ of $\mathcal{G}$ is a **spanning graph** if both $V = W$ and $\mathcal{H}$ is a tree.

Theorem 4.1. Let $\mathcal{G} \equiv (V, E)$ be a graph, then

$$\sum_{v \in V} \deg_{\mathcal{G}}(v) = 2|E| \tag{4.7}$$

Corollary 4.1. For any graph, the number of vertices with odd degree is even.

Definition 4.16. Let $\mathcal{T}$ be a tree, a vertex v is a **leaf** if $\deg_{\mathcal{G}}(v) = 1$.

Proposition 4.1. Every tree with $|V| \geq 2$ has at least two leaves.

Proof. Let $\mathcal{T}$ be a tree. The corollary above suggests that it cannot have one leaf. Consider the case it has no leaf, then since every vertex has at least degree of 2 and $\mathcal{T}$ is connected, there must exist a cycle, which leads to a contradiction. ■

[4]If $\mathcal{G}$ is disconnected, it's impossible for any of its spanning subgraph to be connected.

4.1 Eulerian Graphs

Definition 4.17. Let $\mathcal{G} \equiv (V, E)$ be a graph, then a sequence of vertices $(v_0, v_1, \ldots v_t)$ is an **Eulerian circuit** if

(i) $v_0 = v_t$;

(ii) $v_i v_{i+1} \in E \; \forall i \in \{0, \ldots, t-1\}$;

(iii) $\forall e \in E, \exists \,!i \in \mathbb{Z} \; s.t. \; v_i v_{i+1} = e$.

That's, it is a graph cycle which uses each graph edge exactly once.

Definition 4.18. A graph is **Eulerian** if it contains an eulerian circuit.

Remark 4.2. Some definitions require Eulerian graph to be connected but some don't, check with the lecture notes.

Definition 4.19. A **circuit** is a walk with $x_0 = x_n$.

Theorem 4.2. A graph $\mathcal{G}$ is Eulerian if and only if it is connected and every vertex has even degree.

4.2 Hamiltonian Graphs

Definition 4.20. Let $\mathcal{G} \equiv (V, E)$ be a graph, then a sequence of vertices $(v_0, v_1, \ldots, v_t)$ is a **Hamiltonian cycle** if

1. $v_0 v_t \in E$;
2. $v_i v_{i+1} \in E \; \forall i \in \{0, \ldots, t-1\}$;
3. $\forall v \in V, \exists \; i \in \mathbb{Z} \; s.t. \; v_i = v$.

Definition 4.21. A graph containing Hamiltonian cycle is **Hamiltonian**.

Theorem 4.3. If $\mathcal{G}$ is a graph with n vertices, and $\deg_{\mathcal{G}}(v) \geq \lceil \frac{n}{2} \rceil \; \forall v \in V$, then $\mathcal{G}$ is Hamiltonian.

4.3 Graph Colouring

Definition 4.22. Let $\mathcal{G} \equiv (V, E)$, and C is a set of elements called **colours**. Then a **proper colouring** of $\mathcal{G}$ is a function $\phi : V \to C$ such that

$$\forall x, y \in V, \; xy \in E \implies \phi(x) \neq \phi(y) \tag{4.8}$$

Definition 4.23. The least size of C such that we can construct a proper colouring with it is defined as the **chromatic number** of $\mathcal{G}$, denoted as $\chi(\mathcal{G})$.

Definition 4.24. A graph $\mathcal{G} \equiv (V, E)$ with $\chi(\mathcal{G}) \leq 2$ is called **2-colourable graph**.

Theorem 4.4. A graph is 2-colourable if and only if it does *not* contain an odd cycle.

Proof.
$(\implies)$ Let $\mathcal{G} \equiv (V, E)$ be a 2-clourable graph with proper colouring $\phi : V \to \{\alpha, \beta\}$.
Define $V_1 \equiv \phi^{-1}(\alpha)$ and $V_2 \equiv \phi^{-1}(\beta)$. Clearly those two sets are disjoint and $V = V_1 \cup V_2$.
By definition of proper colouring, for every pair of $x_1, x_2 \in V_1$, $x_1 x_2 \notin E$. The same holds for V_2.
Therefore subgraphs of $\mathcal{G}$ induced from V_1 and V_2 are themselves independent, and $\mathcal{G}$ it bipartite.
We've shown the equivalence between bipartite and 2-colourable.
Suppose there's an odd cycle in $\mathcal{G}$, $C = (x_1, x_2, \ldots, x_n)$, where n is odd.
WLOG, assume $x_1 \in V_1$, by nature of bipartite graph, $x_i \in V_2 \iff i$ even. Therefore $x_n \in V_1$, and for C to be a cycle, we require $x_1 x_n \in E$, which contradicts the fact that $\mathcal{G}$ is bipartite and 2-colourable.
Modus Tollens
$(\impliedby)$ Suppose there exists an odd cycle $C = (x_1, x_2, \ldots, x_n)$ in $\mathcal{G}$, it's easy to show, by induction, that for any proper colouring ϕ of $\mathcal{G}$, $|\phi(C)| \geq 3$. This implies $|\phi(V)| \geq 3$, so $\mathcal{G}$ is not 2-colourable.
Modus Tollens ∎

Definition 4.25. A graph $\mathcal{G} = (V, E)$ is a **bipartite graph** when V can be partitioned into two sets V_1, V_2, such that subgraphs *induced* by V_1 and V_2 are *independent graphs*.

Remark 4.3. Bipartite graphs are 2-colourable. Simply define $\phi : V \to \{\alpha, \beta\}$ as

$$\phi(v) = \alpha \mathbb{1}\{v \in V_1\} + \beta \mathbb{1}\{v \in V_2\} \tag{4.9}$$

Definition 4.26. A **clique** in a graph $\mathcal{G} \equiv (V, E)$ is a set $K \subseteq V$ such that the subgraph induced by K is isomorphic to the $|K|$-complete graph $\mathbf{K}_{|K|}$. (Equivalently, vertices in K are *pair-wise adjacent*)

Definition 4.27. The **maximum clique size** or **clique number** of graph $\mathcal{G}$, denoted as $\omega(\mathcal{G})$ is the largest t such that there exists a clique with t vertices.

Proposition 4.2. For any graph $\mathcal{G}$,

$$\chi(\mathcal{G}) \geq \omega(\mathcal{G}) \tag{4.10}$$

Proposition 4.3 (Generalized Pigeon Hole Principle)**.** Let $f : X \to Y$ be a mapping such that

$$|X| > (m-1)|Y| \tag{4.11}$$

then there exists $\{x_1, \ldots, x_m\} \subseteq X$ such that $f(x_i) = f(x_j)\ \forall i, j$.

Proof. For each $y \in Y$, we can divide X into $|Y|$ partitions, where each partition is defined as the pre-image of one particular $y \in Y$. Let $\{X_i\}$ denote the set of partitions.
We are trying to find the minimum value of $\max_i \{|X_i|\}_{i=1}^{|Y|}$, that's

$$\min_{\text{valid partition}} \max_i \{|X_i|\}_{i=1}^{|Y|} \tag{4.12}$$

the minimum is attained when each partition of X has the same cardinality, which is strictly greater than $m - 1$.
For each of those partitions, it's a pre-image for some value $y \in Y$ with size at least m. ∎

Proposition 4.4. For every $t \geq 3$, there exists a graph $\mathcal{G}_t$ so that $\chi(\mathcal{G}_t) = t$ and $\omega(\mathcal{G}_t) = 2$. *So the difference between χ and ω can be arbitrarily large and this inequality in proposition 4.2 cannot always be tight.*

Definition 4.28. Let $\mathcal{F} = \{S_\alpha : \alpha \in V\}$ be an indexed family of sets, define a graph $\mathcal{G}$ in the following such that

$$S_x \cap S_y \neq \emptyset \iff xy \in E \tag{4.13}$$

Then we call $\mathcal{G}$ an **intersection graph** (representing $\mathcal{F}$).

Remark 4.4. Every graph is an intersection graph since we can explicitly construct a collection of sets from the adjacency matrix of the given graph.

Definition 4.29. $\mathcal{G}$ is an **interval graph** if it is the intersection graph of *a collection of closed intervals in* $\mathbb{R}$.

Theorem 4.5. If $\mathcal{G}$ is an interval graph, then $\chi(\mathcal{G}) = \omega(\mathcal{G})$.

Definition 4.30. A graph $\mathcal{G}$ is **perfect** if $\chi(\mathcal{H}) = \omega(\mathcal{H})$ for every *induced subgraph* $\mathcal{H}$ of $\mathcal{G}$.

Corollary 4.2. Since every induced subgraph of interval graph is also an interval graph, therefore *every interval graph is perfect.*

4.4 Planer Graphs

Definition 4.31. A **drawing** of a graph is a way of associating its vertices with points in $\mathbb{R}^2$ and its edges with simple polygonal arcs whose endpoints are the coordinates associated to the vertices that are the endpoints of the edge.

Definition 4.32. A **planar drawing** of a graph is on in which arcs corresponding to two edges intersect only at a point corresponding to a vertex to which they are both incident.

Definition 4.33. A graph $\mathcal{G}$ is **planar** if it has a planar drawing.

Definition 4.34. A **face** of a *planar drawing* of a graph is a region bounded by edges and vertices and not containing any other vertices or edges.

Theorem 4.6 (Euler's Formula)**.** Let $\mathcal{G}$ be a *connected planer graph* with V vertices and E edges. Then $\mathcal{G}$ has f faces where

$$V - E + f = 2 \tag{4.14}$$

Theorem 4.7. A planar graph with n vertices has at most $3n - 6$ edges when $n \geq 3$.

Theorem 4.8 (Kuratowski's Theorem)**.** A graph is planar <u>if and only if</u> it does not contain either $\mathbf{K}_5$ or $\mathbf{K}_{3,3}$.

Theorem 4.9 (Four Colour Theorem)**.** Every planar graph has chromatic number at most four.

References

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