# MAT246: Concepts in Abstract Mathematics: Lecture 0101 Notes

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# 1 Lecture 1 Sep. 7 2018

**Definition 1.1.** Let  $\mathbb{N} := \{1, 2, 3, ...\}$  be the set of **natural numbers**.

**Theorem 1.1** (Principle of Mathematical Induction). Suppose S is a set of natural numbers,  $S \subseteq \mathbb{N}$ . If

- 1.  $1 \in S$
- 2.  $k \in S \implies k+1 \in S, \forall k \in \mathbb{N}$

then,  $S = \mathbb{N}$ 

**Example 1.1.** Show that

$$1^{2} + 2^{2} + \dots + n^{2} = \frac{n(n+1)(2n+1)}{6} \ \forall n \in \mathbb{N}$$

Proof.

# 2 Lecture 2 Sep. 10 2018

**Theorem 2.1** (Extended Principle of Mathematical Induction). Suppose set  $S \subseteq \mathbb{N}$  and let  $n_0 \in \mathbb{N}$  fixed, if

- 1.  $n_0 \in S$
- 2.  $\forall k \geq n_0, k \in S \implies k+1 \in S$

then  $\{n_0, n_0 + 1, n_0 + 2, \dots\} \subseteq S$ 

**Example 2.1.** Show that

$$n! \ge 3^n \ \forall n \ge 7$$

Proof.

**Theorem 2.2** (Well-Ordering Principle). Every non-empty subset of natural number has a smallest element.

Proof. (Principle of Mathematical Induction)

Let  $S \subseteq \mathbb{N}$ 

Suppose  $1 \in S \land (k \in S \implies k+1 \in S, \forall k \in \mathbb{N})$ 

Show:  $S = \mathbb{N}$ 

Let  $T = \mathbb{N} \backslash S$ 

Suppose  $T \neq \emptyset$ 

By Well-Ordering Principle, there exists a smallest element of T, denoted as  $t_0 \in \mathbb{N}$ . Since  $1 \in S$ , therefore  $t_0 \neq 1$ .

Therefore  $t_0 > 2$ .

Thus  $t_0 - 1 \in \mathbb{N}$  and since  $t_0 = \min T$ ,  $t_0 - 1 \notin T$ 

Therefore  $t_0$  − 1 ∈ S, then,  $t_0$  − 1 + 1 =  $t_0$  ∈ S,

Contradict the assumption that  $t_0 \in T$ .

Thus  $T = \emptyset$  and  $S = \mathbb{N}$ .

**Remark 2.1.** We can use principle of Mathematical Induction to prove Well-Ordering Principle as well.

# 3 Lecture 3 Sep. 12 2018

**Definition 3.1.** Let  $a, b \in \mathbb{N}$  and a divides b, written as a|b if

$$\exists c \in \mathbb{N} \ s.t. \ b = ac$$

And a is a **divisor** of b.

**Definition 3.2.** A natural number p (except 1) is called **prime** if the only divisors of p are 1 and p.

**Lemma 3.1** (Prime numbers are building blocks of natural numbers). Every natural number other than 1 is a  $product^1$  of prime numbers.

**Theorem 3.1** (Principle of Complete Induction). Suppose  $S \subseteq \mathbb{N}$  and if

- 1.  $n_0 \in S$
- 2.  $n_0, n_0 + 1, \dots, k \in S \implies k + 1 \in S, \forall k \ge n_0$

then

$$\{n_0, n_0 + 1, \dots\} \subseteq S$$

*Proof of Lemma*. Let  $S \subseteq \mathbb{N}$  for which the lemma is true,

Want to show:  $S = \mathbb{N} \setminus \{1\}$ 

(Base Case) For 2 it's a product of prime. Thus  $2 \in S$ 

(Inductive Step) Suppose  $\{2, 3, \dots k\} \subseteq S$ 

Consider k + 1, if k + 1 is a prime then k + 1 can be written as a product of itself, as a product of one single prime.

<sup>&</sup>lt;sup>1</sup>Product could mean the product of a single number.

Else, if k + 1 is not a prime, then  $\exists 1 < m, n < k + 1$  s.t. k + 1 = mn.

By induction hypothesis of strong induction, m, n can both be written as product of primes.

 $m = \prod_{i=1}^{\ell} p_i$ ,  $n = \prod_{i=1}^{t} q_i$  where  $p_i$ ,  $q_i$  are all primes. and  $k+1 = \prod_{i=1}^{t} q_i \prod_{i=1}^{\ell} p_i$ 

thus  $k + 1 \in S$ 

by principle of strong induction,  $\{2, 3, \dots, \} \subseteq S$ .

# **Theorem 3.2.** There is no largest prime number.

Proof. (By contradiction)

Assume there is a largest prime p,

then  $\{2, 3, 5, \dots, p\}$  is the set of all primes

Let 
$$M := (2 * 3 * 5 * \cdots * p) + 1 \in \mathbb{N}$$

*M* is either prime or not.

Suppose M is not a prime, then by Lemma 3.1,  $\exists p'$  dividing M.

Obviously  $\forall i \in \{2 * 3 * 5 * \cdots * p\}, i \not\mid M$ .

There is no prime dividing M, which contradict Lemma 3.1

Thus M is a prime, and M > p, which contradicts assumption

Therefore there is no largest prime.

# 4 Lecture 4 Sep. 14 2018

**Theorem 4.1** (the Fundamental Theorem of Arithmetic). Every natural (except 1) is a product of prime(s), and the prime(s) in the product are unique including multiplicity except for the order.

*Proof.* We have already proven that the existential parts of this theorem in Lemma 3.1.

(Proof for the uniqueness part) Suppose there exists natural number (not 1) has 2 different prime factorizations.

By well ordering principle, there is a smallest n, which has two distinct prime factorizations

Say  $n = p_1 p_2 \dots p_k = q_1 q_2 \dots q_\ell$  where  $p_i, q_i$  are all primes.

Notice that  $p_i \neq q_j$  for any combination of (i, j) since if so  $\frac{n}{p_i} = \frac{n}{q_j}$  is a natural number smaller than n having 2 distinct prime factorization, which contradicts our assumption above.

Specifically,  $p_1 \neq q_1$ .

(Case 1:  $p_1 < q_1$ )

Let  $m := n - p_1 q_2 \dots q_\ell \in \mathbb{N}$ 

Notice  $m = p_1(p_2p_3...p_k - q_2q_3...q_\ell)$ 

Also  $m = (q_1 - p_1)(q_2 q_3 \dots q_{\ell})$ 

- $\implies m = p_1 \dots p_k = q_2 q_3 \dots q_\ell (q_1 p_1)$
- $\implies p_1|m$  also notices that  $p_1 \nmid q_2q_3 \dots q_\ell$
- $\implies p_1|(q_1-p_1) \implies p_1|q_1 \implies p_1=q_1$

Contradicts the assumption that  $p_1 < q_1$ 

The other case goes a similar proof.

**Definition 4.1.** A natural number n is called **composite** if it's not 1 or a prime number.

**Remark 4.1.** Natural numbers are partitioned into 3 categories, 1, prime and composite numbers.

**Example 4.1.** Find 20 consecutive composite numbers.

$$(21!) + 2, (21!) + 3, \dots, (21!) + 21$$

**Example 4.2.** Find k consecutive composite numbers.

$$(k+1!)+2,(k+1)!+3,\ldots,(k+1!)+k+1$$

# 5 Lecture 5 Sep. 17 2018

**Definition 5.1.** Let  $a, b \in \mathbb{Z}$ , and let  $m \in \mathbb{N}$ . If m|a-b then we say "a and b are congruent modulo m"

**Remark 5.1.** Regular Induction ← Complete Induction ← Well-Ordering Principle

*Proof.* (WTS: Complete Induction ⇒ Well-Ordering Principle)

Let  $S \subseteq \mathbb{N}$  and  $S \neq \emptyset$ 

(WTS, S has the smallest element)

Assume *S* does not have the smallest element.

Let  $T := S^c$ 

Clearly  $1 \in T$  (prop 1)

Since other wise 1 could be the smallest element of *S*.

Let  $k \in \mathbb{N}$ .

Suppose  $1, 2, 3, ..., k \in T$ , if  $k + 1 \notin T$ , then  $k + 1 \in S$  and k + 1 becomes the smallest element of S and contradicts our assumption above.

Therefore  $1, 2, 3, \dots k \in T \implies k + 1 \in T$ .

By principle of strong induction,  $T = \mathbb{N}$ .

Thus,  $S = \emptyset$ , and contradicts our definition of S.

Therefore  $\forall S \subseteq \mathbb{N} \ s.t. \ S \neq \emptyset$ , S has the smallest element (Well-Ordering Principle).

# **Example 5.1** (Application 2). Is $2^{29} + 3$ divisible by 7?

Solution. Notice  $2^2 \equiv 4 \mod 7$  and  $2^3 \equiv 1 \mod 7$ .

$$\implies (2^3)^9 \equiv 1^9 \mod 7$$

$$\implies 2^{27} \equiv 1 \mod 7$$

$$\implies 2^{29} \equiv 4 \mod 7$$

Also  $3 \equiv 3 \mod 7$ 

$$\implies 2^{29} + 3 \equiv 4 + 3 \mod 7$$

$$\implies 2^{29} + 3 \equiv 7 \mod 7$$

$$\implies 7|2^{29} + 3.$$

**Theorem 5.1** (Rules on computing congruence ). Let  $a, b, c, d \in \mathbb{Z}$  and  $m \in \mathbb{N}$ .

1. 
$$a \equiv b \mod m \land c \equiv d \mod m \implies a + c \equiv b + d \mod m$$

2. 
$$a \equiv b \mod m \land c \equiv d \mod m \implies ac \equiv bd \mod m$$

*Proof.* Let  $a, b, c, d \in \mathbb{Z}$  and  $m \in \mathbb{N}$ ,

suppose  $a \equiv b \mod m \land c \equiv d \mod m$ 

by definition of congruence,  $\exists p, q \in \mathbb{Z} \text{ s.t. } (a-b) = pm \land (c-d) = qm$ 

$$\implies$$
  $(a+c-b-d)=(p+q)m, (p+q)\in\mathbb{Z}$ 

$$\implies a + c \equiv b + d \mod m$$

And  $a = b + pm \wedge c = d + qm$ 

$$ac - bd = (b + pm)(d + qm) - bd$$

$$= bd + dpm + qbm + pqm^2 - bd$$

$$= (dp + qb + pqm)m$$

$$\implies m|ac - bd$$

$$\implies ac \equiv bd \mod m$$

**Proposition 5.1** (Corollary from theorem 5.1).

$$a \equiv b \mod m \implies a + c \equiv b + c \mod m$$

and

$$a \equiv b \mod m \implies a^k \equiv b^k \mod m, \ \forall k \in \mathbb{Z}_{\geq 0}$$

# Lecture 6 Sep. 19 2018

**Theorem 6.1.** Let  $a, b \in \mathbb{Z}$ ,

$$a = b \implies a \equiv b \mod m \ \forall m \in \mathbb{N}$$

**Example 6.1.** What is the reminder when  $3^{202} + 5^9$  is divided by 8

Solution. Notice  $3^2 \equiv 1 \mod 8$ 

Therefore,  $(3^2)^{101} \equiv 1^{101} \mod 8$ 

That's,  $3^{202} \equiv 1 \mod 8$ 

Also  $5^2 \equiv 1 \mod 8$ 

 $\implies (5^2)^4 \equiv 1^4 \mod 8$ 

 $\implies 5^9 \equiv 5 \mod 8$ 

 $\implies$  3<sup>202</sup> + 5<sup>9</sup>  $\equiv$  5 + 1 mod 8

 $\implies$  the reminder is 6.

(Notice that  $3^{202} + 5^9 \equiv 6 \equiv 14 \equiv 22 \equiv \dots \mod 8$ , and the reminder is the smallest integer satisfying above relation.)

**Theorem 6.2.** Let  $M \in \mathbb{Z}$  and  $M = d_N \dots d_2 d_1 d_0, d_i \in \{0, 1, \dots, 9\}^2$ , then

$$3|M\iff 3|\sum_{i=0}^N d_i$$

*Proof.* Notice  $10 \equiv 1 \mod 3$ ,  $100 \equiv 1 \mod 3$  and so on,

(Fact)  $10^k \equiv 1 \mod 3, \ \forall k \in \mathbb{Z}_{\geq 0}$ 

Then  $d_i 10^i \equiv d_i \mod 3$ ,  $\forall i$ Therefore,  $\sum_{i=0}^N 10^i d_i \equiv \sum_{i=0}^N d_i \mod 3$ Therefore  $\sum_{i=0}^N 10^i d_i \equiv 0 \mod 3 \iff \sum_{i=0}^N d_i \equiv 0 \mod 3$ 

**Theorem 6.3.** Let  $M \in \mathbb{Z}$  and  $M = d_N \dots d_2 d_1 d_0, d_i \in \{0, 1, \dots, 9\}$ , then

$$11|M\iff 11|\sum_{i=0}^{N}(-1)^{i}d_{i}$$

*Proof.* Notice  $10^i \equiv (-1)^i \mod 11$ 

Therefore  $10^i d_i \equiv (-1)^i d_i$ 

Thus,  $\sum_{i=0}^{N} 10^{i} d_{i} \equiv \sum_{i=0}^{N} (-1)^{i} d_{i} \mod 11$ Then,  $\sum_{i=0}^{N} 10^{i} d_{i} \equiv 0 \mod 11 \iff \sum_{i=0}^{N} (-1)^{i} d_{i} \equiv 0 \mod 11$ 

<sup>&</sup>lt;sup>2</sup>This means the integer M is constructed from digits  $d_i$ . For example, M = 256, then  $d_0 = 6$ ,  $d_1 = 6$  $5, d_2 = 2$ 

# Lecture 7 Sep. 21 2018

**Theorem 7.1.** Suppose p is a prime and  $a, b \in \mathbb{N}$ , if p|ab then  $p|a \vee p|b$ .

*Proof.* If  $a = 1 \lor b = 1$ , then done. And for the case a = b = 1, the proposition is vacuously true.

Let a, b > 1,

By the fundamental theorem of arithmetic, we can write a, b as their unique prime factorization

$$a = p_1^{\alpha_1} \dots p_k^{\alpha_k}, \ \alpha_i \ge 1 \text{ and } b = q_1^{\beta_1} \dots q_\ell^{\beta_\ell}, \ \beta_i \ge 1$$

then  $a = p_1^{\alpha_1} \dots p_k^{\alpha_k}, \ \alpha_j \ge 1$  and  $b = q_1^{\beta_1} \dots q_\ell^{\beta_\ell}, \ \beta_j \ge 1$ then  $ab = p_1^{\alpha_1} \dots p_k^{\alpha_k} q_1^{\beta_1} \dots q_\ell^{\beta_\ell}$  is the unique prime factorization of ab. Since  $p \in \mathbb{P}$ , therefore,  $p = p_j \lor p = q_j \implies p|a \lor p|b$ 

**Remark 7.1.** We have shown that  $a \equiv b \mod m \implies ca \equiv cb \mod m$ . But notice that

$$ca \equiv cb \mod m \implies a \equiv b \mod m$$

**Definition 7.1.** Let  $a, b \in \mathbb{Z}$ , then we say a and b are **relatively prime** if they have no prime factor in common.

**Theorem 7.2.** Suppose p is a prime and  $a \in \mathbb{Z}$  and  $p \nmid a$ , then  $ax \equiv ay \mod p \implies$  $x \equiv y \mod p$ .

*Proof.* Let  $x, y, a \in \mathbb{N}$  and  $p \in \mathbb{P}$ .

Suppose  $ax \equiv ay \mod p$ 

Then p|a(x - y)

By theorem 7.1,  $p|a \vee p|(x-y)$ 

But by our assumption,  $p \nmid a$ , therefore  $p \mid (x - y)$ 

Thus  $x \equiv y \mod p$ 

**Theorem 7.3** (Generalization of Theorem 7.2). Let  $m \in \mathbb{N}$  and  $a \in \mathbb{Z}$  and a and m are relatively prime. Then

$$ax \equiv ay \mod m \implies x \equiv y \mod m$$

*Proof.* Suppose  $ax \equiv ay \mod m$ 

Then m|a(x - y)

Therefore  $m|a \vee m|(x - y)$ 

For m to divide a, all of m's prime factors have to be in the prime factorization of |a|.

But m and a are relatively prime, therefore  $m \nmid a$ .

Therefore m|(x - y) and that's  $x \equiv y \mod m$ 

**Theorem 7.4.** Any integer a is congruent to mod m to exactly one of  $\{0, 1, \ldots, m-1\}$ .

**Theorem 7.5** (Fermat's Little Theorem). If p is a prime and  $p \nmid a$  (i.e. a and p are relatively prime), then

$$a^{p-1} \equiv 1 \mod p$$

*Proof.* Let  $S := \{a1, a2, \dots a(p-1)\}$ 

Notice that if  $ax_i \equiv ax_i \mod p$ , since  $p \nmid a, x_1 \equiv x_2 \mod p$ .

Since  $1 \le x_i, x_j \le p - 1$ , then  $x_i = x_j$ .

Therefore all elements in S are distinct with mod p

i.e.  $x_i \not\equiv x_i \mod p$ ,  $\forall (i, j) \in \mathbb{Z}^2$ .

Since  $p \not| a \land p \not| m$ ,  $\forall m \in \{1, 2, ..., (p-1)\}$ 

So no element in S is congruent to  $0 \mod p$ .

Thus, S contains p-1 numbers and no two of them are congruent mod p.

Also none of them are congruent to  $0 \mod p$ .

By theorem 7.4, each element in S is congruent to one corresponding element in set  $\{1, 2, \ldots, p-1\}$ .

Therefore  $(a1)(a2)...(a(p-1)) \equiv 1 * 2 * ... * (p-1) \mod p$ 

That's  $a^{p-1}(1*2*\cdots*(p-1)) \equiv 1*2*\cdots*(p-1) \mod p$ 

Clearly  $p \nmid (1 * 2 * ... (p-1))$ , since if a prime divides a product of natural numbers, the prime must divide at least one of elements in the product.

Therefore  $a^{p-1} \equiv 1 \mod p$ 

# 8 Lecture 8 Sep. 24 2018

**Definition 8.1.** Let  $p \in \mathbb{N}$  and  $a \in \mathbb{Z}$ . The **multiplicative inverse** mod p of a is an integer b such that

$$ab \equiv 1 \mod p$$

**Remark 8.1.** Notice that the multiplicative inverse is generally not unique but unique up to  $\mod p$ .

**Corollary 8.1.** Let  $p \in \mathbb{P}$ ,  $a \in \mathbb{N}$  and  $p \nmid a$ . Then

$$\exists b \in \mathbb{Z}, \ s.t. \ ba \equiv 1 \mod p$$

*Proof.* Let  $p \in \mathbb{Z}$  and  $a \in \mathbb{Z}$ Suppose  $p \nmid a$ , then by Fermat's little theorem,  $a^{p-1} \equiv 1 \mod p \implies a^{p-2}a \equiv 1 \mod p$ Take  $b = a^{p-2} \in \mathbb{Z}$  and  $ab \equiv 1 \mod p$ 

**Example 8.1.** Let a = 8 and p = 5. Obviously  $p \nmid a$ . By corollary above,

$$\exists b \in \mathbb{Z}, \ s.t. \ 8b \equiv 1 \mod 5$$

Notice b = 2 satisfies above equation.

**Remark 8.2.** Corollary 8.1 requires *p* to be a prime.

**Corollary 8.2** (Generalization). Let a and  $m \in \mathbb{N}$  and a and m are relatively prime, then

$$\exists b \in \mathbb{Z}, \ s.t. \ ab \equiv 1 \mod m$$

**Theorem 8.1** (Wilsons' Theorem). Let  $p \in \mathbb{P}$  then

$$(p-1)! \equiv -1 \mod p$$

*Proof.* Let  $p \in \mathbb{P}$ 

if  $p = 2 \lor p = 3$ , then  $1! \equiv -1 \mod 2$  and  $2! \equiv -1 \mod 3$ .

Otherwise, suppose p > 3,

Consider, let  $S := \{2, 3, 4, \dots, p-2\}$ 

Notice that none of S is divisible by p.

Therefore p is relatively prime to all elements in S.

Then by Corollary 8.1,  $\exists b_i \in \mathbb{Z} \ s.t. \ b_i s_i \equiv 1 \mod p, \ \forall s_i \in S$ .

Notice that 0 has no multiplicative inverse and

$$(p-1)(p-1) = p^2 - 2p + 1 \equiv 1 \mod p$$

That's, 1 and (p-1) have themselves as their multiplicative inverse.

Also notice that for any  $s_i \in S$ ,  $s_i$  does not have itself as its multiplicative inverse.

If  $a \in S$  has itself as it's multiplicative inverse, then

$$a^{2} \equiv 1 \mod p$$

$$\implies a^{2} - 1 \equiv 0 \mod p$$

$$\implies (a+1)(a-1) \equiv 0 \mod p$$

$$\implies p|(a+1)(a-1)$$

Notice that at last one of (a + 1) and (a - 1) is in set S since  $p > 3 \implies S \neq \emptyset$ . This contradicts what we argued above, *none of* S *is divisible by* p. That's

$$s_i s_i \not\equiv 1 \mod p, \ \forall s_i \in S$$

Note that if y is a multiplicative inverse of x, then x is a multiplicative inverse of y. Notice that for any  $s_i \in S$ , by Corollary 8.1,

there exists an integer  $b_i$  s.t.  $s_i b_i \equiv 1 \mod p$ 

And the multiplicative inverse is unique up to  $\mod p$ ,

Thus  $s_i(b_i \mod p) \equiv 1 \mod p$  and  $(b_i \mod p) \in S$ .

And for all elements in S has one of their multiplicative inverse in S,

That's

$$s_i s_j \equiv 1 \mod p, \ i \neq j$$

Notice p > 3 implies p is odd, so |S| is even.

Match every pair of multiplicative inverses in S and they collapse to  $1 \mod p$  Therefore

$$2 \cdot 3 \cdot 4 \cdots (p-2) \equiv 1 \mod p$$

$$\implies 2 \cdot 3 \cdot 4 \cdots (p-2) \cdot (p-1) \equiv (p-1) \mod p$$

$$\implies (p-1)! \equiv -1 \mod p$$

# 9 Lecture 9 Sep. 26 2018

**Remark 9.1.** Recall that an integer n is even iff  $n \equiv 0 \mod 2$  and is odd iff  $n \equiv 1 \mod 2$ .

**Theorem 9.1.** There are infinitely many primes of the form 4k + 3, where  $k \in \mathbb{Z}$ .

*Proof.* Note that odd numbers n can be classified as  $n \equiv 1 \mod 4$  and  $n \equiv 3 \equiv -1 \mod 4$ 

(Suppose 1) there are only finitely many primes in the form 4k + 3.

Let finite set  $S := \{p_1, p_2, \dots p_m\}$  denotes the collection of them.

And notice that  $p_i \equiv -1 \mod 4$ ,  $\forall p_i \in S$ .

Let

$$M:=(p_1\cdot p_2\cdots p_m)^2+2$$

and  $M \equiv 1 + 2 \equiv 3 \equiv -1 \mod 4$ .

Therefore *M* is an odd natural number.

By the Fundamental Theorem of Arithmetic, M can be factorized into product of

primes.

$$M = \prod_{i=1}^{\ell} q_i$$

and since M is odd,  $q_i \neq 2 \ \forall i$ . Thus all  $q_i$  are odd.

(Suppose 2) All  $q_i \equiv 1 \mod 4$ .

Then  $M \equiv 1 \mod 4$ .

Contradict the fact that  $M \equiv -1 \mod 4$ . Thus (Suppose 2) is false.

Therefore  $\exists i, s.t. q_i \equiv -1 \mod 4$ .

From (Suppose 1), S is the collection of all primes that  $\equiv -1 \mod 4$ .

Therefore  $q_i = p_j$  for some j.

Therefore  $p_i|M$ .

Also note that  $p_i|(p_1 \cdot p_2 \cdots p_m) \implies p_i|(p_1 \cdot p_2 \cdots p_m)^2$ 

 $\implies p_i|2 \implies p_i = 2$  contradicts the fact that  $p_i$  is odd.

Therefore (Suppose 1) is false, there are infinitely many primes taking the form 4k + 3.

# **Example 9.1.** Find $7^{20^{30}} \mod 5$ .

Solution. Let  $n := 20^{30}$ .

Notice that  $7^4 \equiv 1 \mod 5$ .

And if  $n \equiv r \mod 4$  where  $r \in \mathbb{Z}$ ,

n = 4k + r and  $7^n \equiv 7^{4k+r} \equiv (7^4)^k \times 7^r \equiv 1^k \times 7^r \equiv 7^r \mod 5$ .

Notice that  $20 \equiv 0 \mod 4 \implies 20^{30} \equiv 0 \mod 4$ .

Thus r = 0.

Therefore  $7^n \equiv 7^0 \equiv 1 \mod 5$ .

Thus  $7^{20^{30}} \mod 5 = 1$ .

# **Example 9.2.** Find $10^{3^{30}} \mod 7$ .

Solution. Notice that  $10^6 \equiv 1 \mod 7$ .

And  $3 \equiv 3 \mod 6$ ,  $3^2 \equiv 3 \mod 6$ ,  $3^3 \equiv 3 \mod 6$ ...

Using induction, we can show that

$$3^k \equiv 3 \mod 6, \ \forall k \in \mathbb{Z}_{\geq 0}$$

Therefore  $3^{30} \equiv 3 \mod 6$ .

That's  $3^{30} = 6k + 3$  for some *k*.

Thus  $10^{3^{30}} \equiv (10^6)^k \times 10^3 \equiv (1)^k \times 10^3 \equiv -1 \equiv 6 \mod 7$ . So  $10^{3^{30}} \mod 7 = 6$ .

#### Lecture 10 Sep. 28 2018 10

**Example 10.1.** Find  $8^{9^{10^{11}}}$ mod 5.

*Solution.* Let  $n := 9^{10^{11}}$ 

And notices that  $8^4 \equiv 1 \mod 5$ .

Then find  $n \mod 4$ 

Note that  $9 \equiv 1 \mod 4 \implies 9^{10^{11}} \equiv 1 \mod 4$ .

Thus n = 4k + 1. Therefore  $8^{9^{10^{11}}} \equiv (8^4)^k \cdot 8 \equiv 1 \cdot 3 \mod 5$ . That's  $8^{9^{10^{11}}} \mod 5 = 3$ .

**Definition 10.1** (Euler  $\phi$ -function). Let  $m \in \mathbb{N}$  and  $\phi(m) : \mathbb{N} \to \mathbb{N}$  is defined as the number of elements in  $\{1, 2, ..., m-1\}$  that are relatively prime to m.

**Example 10.2.** For m = 8, note that  $\{1, 3, 5, 7\} \subset \{1, 2, \dots, 7\}$  are relatively prime with 8, therefore  $\phi(8) = 4$ .

And for m = 11, since m is a prime, then every integer between 1 and m - 1 are relatively prime with 11. Therefore  $\phi(11) = 10$ .

And notice that  $\phi(p) = p - 1$  if  $p \in \mathbb{P}$ . (Fermat's Little Theorem)

**Proposition 10.1.** Let p, q be two distinct primes, then

$$\phi(pq) = (p-1)(q-1)$$

*Proof.* Let  $S := \{1, 2, ..., pq - 1\}.$ 

WLOG, assume p < q.

We need find all elements in S that with either p or q in their prime factorization to find elements in S that are not relatively prime to pq.

And those elements are multiples of p and multiples of q.

And since  $pq \notin S$ , the largest multiple of p in S is (q-1)p and the largest multiple of q in S is q(p-1).

And since there is no multiple of both p and q in set S, therefore there's no overlapping between multiples of p and multiples of q.

Therefore exists (p-1) + (q-1) elements that are not relatively. prime to pq.

Therefore  $\phi(pq) = (pq - 1) - (p - 1) - (q - 1)$ 

$$= pq - p - q + 1$$

$$= (p-1)(q-1)$$

**Proposition 10.2.** For any natural number  $m \in \mathbb{N}$ . Therefore m can be expressed as

$$m = p_1^{\alpha_1} p_2^{\alpha_2} \cdots p_k^{\alpha_k}$$

Then

$$\phi(m) = \phi(p_1^{\alpha_1})\phi(p_2^{\alpha_2})\cdots\phi(p_k^{\alpha_k})$$

And

$$\phi(p^{\alpha}) = p^{\alpha} - p^{\alpha - 1} = p^{\alpha - 1}(p - 1)$$

Therefore

$$\phi(m) = (p_1^{\alpha_1} - p_1^{\alpha_1 - 1})(p_2^{\alpha_2} - p_2^{\alpha_2 - 1}) \cdots (p_k^{\alpha_k} - p_k^{\alpha_k - 1})$$

#### Example 10.3.

$$\phi(6) = \phi(2^{1}3^{1})$$

$$= \phi(2^{1})\phi(3^{1})$$

$$= (2^{1} - 2^{0})(3^{1} - 3^{0})$$

$$= (2 - 1)(3 - 1) = 2$$

#### Example 10.4.

$$\phi(8) = \phi(2^3)$$
$$= (2^3 - 2^2) = 4$$

**Theorem 10.1** (Euler's Theorem). Suppose  $m \in \mathbb{N} \setminus \{1\}$ . And  $a \in \mathbb{N}$  <sup>3</sup>Assume a and m are relatively prime, then

$$a^{\phi(m)} \equiv 1 \mod m$$

**Remark 10.1.** This theorem is a generalization of Fermat's Little Theorem. When  $m \in \mathbb{P}$ , it becomes Fermat's Little Theorem.

*Proof.* Let  $S := \{r_1, r_2, \dots r_{\phi(m)}\}$  be the set of all elements in  $\{1, 2, \dots, m-1\}$  that are relatively prime to m.

Let 
$$T := \{ar_1, ar_2, \dots ar_{\phi(m)}\}.$$

(Observation 1) that no two elements in S are congruent to each other  $\mod m$ . Since all elements are in the range [1, m-1] and they are the reminder while  $r_i$  is divided by m.

<sup>&</sup>lt;sup>3</sup>Also true for  $a \in \mathbb{Z}$ 

Also notice that elements in T are not congruent to each other  $\mod m$ . Since, suppose

$$ar_i \equiv ar_i \mod m$$

for some (i, j).

Since a and m are relatively prime, therefore we could use cancellation law.

$$r_i = \equiv r_j \mod m$$

This would contradict our observation 1

(Observation 2) elements in T are not congruent to each other  $\mod m$ .

Therefore elements in S are congruent to elements in  $T \mod m$  in some order. Therefore

$$r_1 r_2 r_3 \cdots r_{\phi(m)} \equiv a^{\phi(m)} r_1 r_2 \cdots r_{\phi(m)} \mod m$$

And notice  $r_1r_2r_3\cdots r_{\phi(m)}$  is a product of natural numbers relatively prime to m. Therefore  $r_1r_2r_3\cdots r_{\phi(m)}$  is relatively prime to m. And by cancellation law, we have

$$a^{\phi(m)} \equiv 1 \mod m$$

# 11 Lecture 11 Oct. 1 2018

#### 11.1 Rational and Irrational Numbers

**Definition 11.1.** A rational number is an expression in form

$$\frac{m}{n}$$
,  $m, n \in \mathbb{Z}$ ,  $n \neq 0$ 

**Definition 11.2.** Two rational numbers  $\frac{m_1}{n_1}$ ,  $\frac{m_2}{n_2} \in \mathbb{Q}$  are **equal** if and only if  $m_1 n_2 = m_2 n_1$ .

**Definition 11.3.** Arithmetic on  $\mathbb{Q}$  are defined as

- Addition + :  $\frac{m_1}{n_1} + \frac{m_2}{n_2} := \frac{m_1 n_2 + m_2 n_1}{n_1 n_2}$
- Multiplication  $\times$  :  $\frac{m_1}{n_1} \times \frac{m_2}{n_2} := \frac{m_1 m_2}{n_1 n_2}$
- **Subtraction**  $-: \frac{m_1}{n_1} \frac{m_2}{n_2} := \frac{m_1 n_2 m_2 n_1}{n_1 n_2}$

• **Division**  $\div$  :  $\frac{\frac{m_1}{n_1}}{\frac{m_2}{n_2}}$  :=  $\frac{m_1 n_2}{n_1 m_2}$ , defined only if  $m_2 \neq 0$ .

**Definition 11.4.** The **multiplicative inverse** of a <u>non-zero</u> rational number  $x \ne 0$  is a rational number y such that xy = 1.

**Remark 11.1.** Let  $x = \frac{m}{n} \neq 0$ , then the multiplicative inverse  $y = \frac{n}{m}$ .

**Example 11.1.** Claim:  $\sqrt{2}$  is not rational.

*Proof.* Assume  $\sqrt{2}$  is rational,

by definition of rational numbers,  $\sqrt{2} = \frac{m}{n}$  where  $m, n \in \mathbb{Z}, n \neq 0$ .

Divide numerator and denominator by their common prime factors (if any).

Assume m and n have been reduced so that they are relatively prime.

$$\implies 2 = \frac{m^2}{n^2}$$

$$\iff 2n^2 = m^2$$

$$\implies 2|m^2$$

Consider if  $2 \nmid m$ , then m is odd, then  $2 \nmid m^2$ . Take the contraposition,  $2|m^2 \implies 2|m$ .

$$\implies 2|m$$

$$\implies m = 2q, \ q \in \mathbb{Z}$$

$$\implies 2n^2 = 4q^2$$

$$\implies n^2 = 2q^2$$

$$\implies 2|n^2$$

$$\implies 2|n$$

That's  $2|m \wedge 2|n$ , which contradicts our assumption that m and n are relatively prime. Therefore  $\sqrt{2}$  cannot be rational.

**Definition 11.5** (non-rigorous definition). **Real numbers**, denoted as  $\mathbb{R}$ , are numbers representing distance of points on a line from 0.

**Definition 11.6. Irrational numbers** are real numbers which are not rational.  $(\mathbb{R}\backslash\mathbb{Q})$ 

**Proposition 11.1.** Let  $p \in \mathbb{P}$  and  $m \in \mathbb{Z}$ , then

$$p|m^2 \implies p|m$$

*Proof.* Let  $m = q_1 q_2 \dots q_\ell$  be the unique prime factorization.

Suppose  $p \nmid m$ , then  $p \notin \{q_1, q_2, \dots, q_\ell\}$ . Obviously,  $m^2 = q_1^2 q_2^2 \dots q_\ell^2$  as it's prime factorization. Then  $p \nmid m^2$ .

**Example 11.2.**  $\sqrt{p} \notin \mathbb{Q}, \ \forall p \in \mathbb{P}.$ 

*Proof.* Let  $p \in \mathbb{P}$ , Suppose  $\sqrt{p} \in \mathbb{Q}$ .

Therefore  $\sqrt{p} = \frac{m}{n}$  where  $m, n \in \mathbb{Z}$  and  $n \neq 0$ .

Assume  $\frac{m}{n}$  has been reduced such that m and n are relatively prime.

$$\implies pn^2 = m^2$$

$$\implies p|m^2$$

$$\implies p|m$$

$$\implies m = pr, \ r \in \mathbb{Z}.$$

$$\implies pn^2 = p^2r^2$$

$$\implies n^2 = pr^2$$

$$\implies p|n^2$$

$$\implies p|n$$

$$\implies p|n$$

Contradicts the assumption that m and n are relatively prime.

#### 12 Lecture 12 Oct. 3 2018

**Definition 12.1.** A natural number (other than 1) is called a **perfect square** if it is the square of some natural number.

**Theorem 12.1.** A natural number m is a perfect square if and only if every prime factor occurs with an even power in its prime decomposition.

*Proof.* ( $\Longrightarrow$ ) Suppose m is a perfect square,

Then  $m = n^2, b \in \mathbb{N}$ .

Let  $n = p_1^{\alpha_1} \dots p_k^{\alpha_k}$  be the prime decomposition. Then  $m = p_1^{2\alpha_1} \dots p_k^{2\alpha_k}$ .

Obviously all prime factors in the prime factorization occurs with an even power.

( $\iff$ ) Suppose  $m=p_1^{2\alpha_1}\dots p_k^{2\alpha_k}$  as its prime decomposition. Then  $m=(p_1^{\alpha_1}\dots p_k^{\alpha_k})^2$  and  $n=p_1^{\alpha_1}\dots p_k^{\alpha_k}\in\mathbb{N}$ .

Therefore *m* is a perfect square.

**Theorem 12.2** (Generalization). Let  $n \in \mathbb{N}$  other than 1, then <sup>4</sup>

$$\sqrt{n} \in \mathbb{Q} \iff n \text{ is a perfect square}$$

*Proof.* ( $\iff$ ) if *n* is perfect square, then  $\sqrt{n} \in \mathbb{N}$ .

Obviously a natural number is rational.

 $(\Longrightarrow)$  Suppose  $\sqrt{n} \in \mathbb{Q}$ .

Then

$$\sqrt{n} = \frac{m}{l} \in \mathbb{Q}$$

where  $m, l \in \mathbb{Z}$  and  $l \neq 0$ .

Since  $\sqrt{n} > 0$ , WLOG, assume  $m, l \ge 0$ .

Suppose m, l are relatively prime. (Otherwise, factorize the friction so that m and l are relatively prime.)

Then

$$m^2 = nl^2$$

(Suppose 1) l > 1 and p is a prime in the prime decomposition of m, i.e. p|l. Thus  $p|l^2$  and therefore  $p|m^2$ .

By proposition 11.1 (previous lecture), p|m

And we have  $p|l \wedge p|m$  which contradicts our assumption that m, l are relatively prime.

Therefore (Suppose 1) is false and  $l \le 1$  (so that l has no prime factor).

Also notice that  $l \in \mathbb{Z}$  and  $l \ge 0$ . therefore l = 1.

Therefore  $n = m^2$  and n is a prefect square.

# **Example 12.1.** Claim $\sqrt[3]{4}$ is irrational.

*Proof.* Suppose  $\sqrt[3]{4}$  is rational and

$$\sqrt[3]{4} = \frac{m}{n} \implies 4 = \frac{m^3}{n^3} \implies 2^2 n^3 = m^3$$

Suppose

$$n = p_1^{\alpha_1} \dots p_k^{\alpha_k}$$
$$m = q_1^{\beta_1} \dots q_\ell^{\beta_\ell}$$

<sup>&</sup>lt;sup>4</sup>The square root here denotes the positive square root.

The prime factor 2 has power of 2 or  $2 + 3\alpha_i$  on the left hand side.

And have power of  $3\beta_i$  on the right hand side.

The left hand side power is congruent to 2 mod 3 and the right hand side is congruent to 0 mod 3.

It's impossible for them to be equal. Thus, contradicts the uniqueness of prime decomposition.

Therefore  $\sqrt[3]{4}$  cannot be rational.

#### Lecture 13 Oct. 5 2018 13

**Example 13.1.**  $\sqrt{3} + \sqrt{5}$  is irrational.

*Proof.* Suppose  $\sqrt{3} + \sqrt{5}$  are rational then  $\sqrt{3} + \sqrt{5} = \frac{m}{n}$ .

$$\implies \sqrt{5} = \frac{m}{n} - \sqrt{3}$$

$$\implies 5 = (\frac{m}{n} - \sqrt{3})^2 = \frac{m^2}{n^2} - \frac{2m\sqrt{3}}{n} + 3$$

$$\implies \sqrt{3} = \frac{5 - 3 - \frac{m^2}{n^2}}{-\frac{2n}{n}}$$

Obviously the right hand side is rational, leads to contradiction.

Therefore  $\sqrt{3} + \sqrt{5} \notin \mathbb{Q}$ .

**Example 13.2.** Are there two irrational numbers x, y such that  $x^y \in \mathbb{Q}$ ?

Solution. Consider  $\sqrt{3}\sqrt{2}$ .

case 1:  $\sqrt{3}^{\sqrt{2}} \in \mathbb{Q}$ , then take  $x = \sqrt{3}$  and  $y = \sqrt{2}$ .

**case 2:**  $\sqrt{3}^{\sqrt{2}} \notin \mathbb{Q}$ , then take  $x = \sqrt{3}^{\sqrt{2}}$ ,  $y = \sqrt{2}$ . And  $x^y = \sqrt{3}^{\sqrt{2}\sqrt{2}} = \sqrt{3}^2 = 3 \in \mathbb{Q}$ .

**Remark 13.1.** Basic arithmetic operations on integers preserve rationality.

**Theorem 13.1** (Rational Root Theorem). Consider a polynomial with integer coefficients,

$$a_0 + a_1 x + x_2 x^2 + \dots + a_k x^k, \ a_i \in \mathbb{Z}$$

If  $\frac{m}{n}$  is a rational root for the polynomial and m and n are relatively prime. Then

$$m|a_0 \wedge n|a_k$$

*Proof.* Suppose  $\frac{m}{n}$  is a rational root for the polynomial, then

$$a_0 + a_1 \frac{m}{n} + a_2 \frac{m^2}{n^2} + \dots + a_k \frac{m^k}{n^k} = 0$$

Thus

$$a_0 n^k + a_1 m n^{k-1} + a_2 m^2 n^{k-2} + \dots + a_k m^k = 0$$

And

$$-a_0n^k = a_1mn^{k-1} + a_2m^2n^{k-2} + \dots + a_km^k$$

Therefore  $m|a_0n^k$  and since  $m \nmid n$ , thus  $m|a_0$ . Similarly,

$$-a_k m^k = a_0 n^k + a_1 m n^{k-1} + a_2 m^2 n^{k-2} + \dots + a_{k-1} m^{k-1} n$$

Which implies  $n|a_km^k$  and since  $n \nmid m$ , thus  $n|a_k$ .

# 14 Lecture 14 Oct. 10 2018 Euclidean Algorithm

**Definition 14.1.** The **greatest common divisor**(gcd) of  $m, n \in \mathbb{N}$  is denoted as gcd(m, n) or (m; n) is the largest natural number that divides both m and n.

Example 14.1.

$$gcd(27, 15) = 3$$
  
 $gcd(36, 48) = 12$   
 $gcd(7, 21) = 7$ 

#### 14.1 Conventional Method

Factorize m and n into primes and in general,

$$m = p_1^{\alpha_1} \dots p_k^{\alpha_k}$$
$$n = p_1^{\beta_1} \dots p_k^{\beta_k}$$

where  $\{p_1, \dots p_k\} = \{\text{prime factors of } m\} \cup \{\text{ prime factors of } n\}$ . And  $\alpha_i, \beta_i \ge 0$ . and gcd could be found by

$$gcd(m,n) = p_1^{\min\{\alpha_1,\beta_1\}} \dots p_k^{\min\{\alpha_k,\beta_k\}}$$

## 14.2 Euclidean Algorithm

Notice that  $r|a, b \implies r|gcd(a, b)$ . For  $a, b \in \mathbb{N}$ , WLOG, assuming  $a \ge b$ .

$$a = q_0b + r_0 \quad r_0, q_0 \in \mathbb{N}, \ 0 \le r_0 < b$$

$$b = q_1r_0 + r_1 \quad r_1, q_1 \in \mathbb{N}, \ 0 \le r_1 < r_0$$

$$r_0 = q_2r_1 + r_2$$

 $r_i$  is strictly decreasing, and it's guaranteed to be 0 after certain iterations.

 $r_{k-2} = q_k r_{k-1} + r_k$  $r_{k-1} = q_{k+1} r_k + 0$ 

and then  $r_k$  is the greatest common divisor of a and b.

*Proof.* WTS:  $r_k = gcd(a, b)$ 

Obviously  $r_k | r_{k-1}$ 

Then,  $r_k|r_{k-1} \wedge r_k|r_k \implies r_k|r_{k-2}$ 

Similarly, tracing upwards through the Euclidean Algorithm,

we have  $r_k|b$  and  $r_k|a$ .

so  $r_k \le gcd(a, b)$  since  $r_k$  is a common divisor of a and b.

Since  $gcd|a \wedge gcd|b$ ,

Therefore  $gcd|r_0$ ,

similarly, tracing downloads in the Euclidean Algorithm,

 $gcd|r_k$ , so  $gcd \leq r_k$ .

Therefore  $r_k = gcd(a, b)$ 

**Theorem 14.1.** Given natural numbers a and b with the greatest common divisor d, there exists integers x and y such that

$$d = ax + by$$

*Proof.* This can be seen by working upwards in the sequence of equations that constitute the Euclidean Algorithm.

# 15 Lecture 15 Oct. 12 2018 Public Key Cryptography, RSA Public Key

**Lemma 15.1.** Let N = pq where  $p \neq q$  are distinct primes, and let n and M be integers.

Then

$$n \equiv 1 \mod \phi(N) \implies M^n \equiv M \mod N$$

*Proof.* Note that  $\phi(N) = \phi(pq) = (p-1)(q-1)$ .

And suppose  $\phi(N)|(n-1)$ ,

Then  $k\phi(N) = n - 1$  for some k.

That's  $n = 1 + k\phi(N)$ .

Therefore  $M^n = M^{1+k\phi(N)} = (M^{\phi(N)})^k \cdot M$ 

It's sufficient to show  $M^n \equiv M \mod N$ 

by showing  $M^n \equiv M \mod p$  and  $M^n \equiv M \mod q$ 

To show  $M^n \equiv M \mod p$ ,

Case 1: p|M, then  $0^n \equiv 0 \mod p$ , done.

Case 2:  $p \not\mid M$ , then  $M^n = (M^{\phi(N)})^k \cdot M = (M^{(p-1)})^{(q-1)k} \cdot M$ 

By Fermat's Little Theorem,  $M^{p-1} \equiv 1 \mod p$ .

Therefore  $M^n \equiv 1^{(q-1)k} \times M \mod p$ .

# 15.1 RSA Public Key Procedures

Procedures:

- 1. **Receiver:** pick two large distinct primes  $p \neq q$  and calculate  $N = p \times q$ .
- 2. **Receiver:** calculate  $\phi(N) = (p-1)(q-1)$  and pick *e* relatively prime to  $\phi(N)$ .
- 3. **Receiver:** announce N and e.
- 4. **Sender:** choose message  $M \in \mathbb{N}$  satisfies M < N (if  $M \ge N$ , break M into pieces.)
- 5. **Sender:** find  $M^e \equiv R \mod N$ .
- 6. **Sender:** announce the encoded message *R*.
- 7. **Receiver:** pick  $d \ge 0$  s.t.  $de + k\phi(N) = 1$  as the decoder. Such z-linear combination is guaranteed to exist.
- 8. **Receiver:** the original message M can be found by  $R^d \equiv M \mod N$

$$Proof. \ R^d \equiv (M^e)^d \equiv M^{ed} \mod N$$

Since  $ed \equiv 1 \mod \phi(N)$ 

By lemma 15.1,  $M^{ed} \equiv M \mod N$ .

# 16 Lecture 16 Oct. 15 2018 RSA Cryptography Examples

### 16.1 Recall

- 1. **Receiver:** choose  $p, q \in \mathbb{P}$  and computes  $N = pq, \phi(N) = (p-1)(q-1)$  and choose e *s.t.*  $gcd(e, \phi(N)) = 1$ . Then announces e, N.
- 2. **Sender:** choose  $0 \le M < N$  and calculate R such that  $R = M^e \mod N$ .
- 3. **Receiver:** compute decoder d s.t.  $de + k\phi(N) = 1$ . And decode message  $M^* = R^d \mod N$

## **16.2** More Examples

**Example 16.1. Receiver:** pick p = 11, q = 7. Calculate N = 77 and  $\phi(N) = 10 * 6 = 60$ . Pick e = 13 which is relatively prime to  $\phi(N)$ .

**Receiver:** announces N = 77 and e = 13 to sender.

**Sender:** pick message M = 71 < N and *encodes* message by computing  $71^{13} \equiv R \mod 77$ .

$$71 \equiv -6 \mod 77$$

$$71^{3} \equiv (-6)^{3} \equiv 216 \equiv 15 \mod 77$$

$$(71)^{6} \equiv (71^{3})^{2} \equiv 15^{2} \equiv 225 \equiv -6 \mod 77$$

$$(71)^{12} \equiv (71^{6})^{2} \equiv (-6)^{2} \equiv 36 \mod 77$$

$$(71)^{13} \equiv 36 \times (-6) \equiv -216 \equiv 15 \mod 77$$

And calculate R = 15 satisfies  $71^{13} \equiv 15 \mod 77$ .

**Sender:** announces R = 15 to the rest of world.

**Receiver:** find  $d \ge 0$  satisfying  $d \times e + k \times \phi(N) = 1$ . as the *decoder*. And find

that d = 37, k = -8.

**Receiver:** compute  $R^d \mod 77$ 

$$15^{2} \equiv 225 \equiv -6 \mod 77$$

$$15^{6} \equiv -216 \equiv 15 \mod 77$$

$$15^{12} \equiv 15^{2} \equiv -6 \mod 77$$

$$15^{24} \equiv 36 \mod 77$$

$$15^{36} \equiv -216 \equiv 15 \mod 77$$

$$15^{37} \equiv 15^{2} \equiv 226 \equiv -6 \equiv 71 \mod 77$$

$$\implies M^{*} = 15^{37} \mod 77 = 71$$

**Security of RSA** For anyone knowing N but does not know  $\phi(N)$ . To compute the decoder d,  $\phi(N)$  needs to be calculated.

- 1. **Method 1:** Use definition and iterating through  $\{1, 2, ... N\}$  and compute  $\phi(N)$ .
- 2. **Method 2:** Factorize N and find p and q, then calculate  $\phi(N) = (p-1)(q-1)$ .

Both brute force methods are impractical in terms of run-time.

**Example 16.2. Receiver:** pick p = 11, q = 7 then N = pq = 77 and  $\phi(N) = (p-1)(q-1) = 60$  and choose e = 13.

**Receiver:** announce N = 77 and e = 13.

**Sender:** pick M = 76 < 77 and  $M^{13} \equiv (-1)^1 3 \equiv -1 \equiv 76 \mod 77$ . Announce R = 76.

**Receiver:** find *decoder*  $d \ge 0$  s.t.  $d \times 13 + k \times 60 = 1$ . Found d = 13.

**Receiver:** Compute  $R^d \mod 77$ .

$$R = 76 \equiv -1 \mod 77$$

$$R^{13} \equiv (-1)^{13} \equiv -1 \equiv 76 \mod 77$$

$$\implies M^* = R^{13} \mod 77 = 76$$

#### 17 Lecture 17 Oct. 17 2018

**Remark 17.1.** In RSA, picking the *decoder*  $d \ge 0$  s.t.  $de + \phi(N)k = 1$  is equivalent to pick d such that

$$de \equiv 1 \mod \phi(N)$$

#### 17.1 Chinese Reminder Theorem

**Theorem 17.1** (Chinese Reminder Theorem (CRT)). Solve system of congruent equations, where  $m_1$  and  $m_2$  are relatively prime,

$$\begin{cases} x \equiv a \mod m_1 \\ x \equiv b \mod m_2 \end{cases}$$

The solution is given by

$$x = ax_2m_2 + bx_1m_1$$

where  $x_1$  and  $x_2$  satisfy

$$\begin{cases} x_1 m_1 \equiv 1 \mod m_2 \\ x_2 m_2 \equiv 1 \mod m_1 \end{cases}$$

The general solution x is the superposition of two specific solutions.

*Proof.* If  $m_1$  and  $m_2$  are relatively prime, by Theorem 14.1,

$$\exists x_1, x_2 \in \mathbb{Z} \ s.t. \ x_1m_1 + x_2m_2 = 1$$

Taking congruence with respect to mod  $m_1$  and  $m_2$  gives

$$\begin{cases} 1 \equiv x_2 m_2 \mod m_1 \\ 1 \equiv x_1 m_1 \mod m_2 \end{cases}$$

Consider

$$x = ax_2m_2 + bx_1m_1$$

Clearly

$$x - a = a(x_2m_2 - 1) + bx_1m_1$$

$$m_1|(x_2m_2 - 1) \land m_1|bx_1m_1 \implies m_1|x - a$$

$$\implies x \equiv a \mod m_1$$

Similarly, we can show  $x \equiv b \mod m_2$ .

Thus *x* is the solution to system of equations

$$\begin{cases} x \equiv a \mod m_1 \\ x \equiv b \mod m_2 \end{cases}$$

Example 17.1. Solve

$$\begin{cases} x \equiv 5 \mod 7 \\ x \equiv 13 \mod 8 \end{cases}$$

Solution. Solve

$$\begin{cases} x_1 \times 7 \equiv 1 \mod 8 \\ x_2 \times 8 \equiv 1 \mod 7 \end{cases}$$

Solve  $x_1 = 7$  and  $x_2 = 1$ . And one solution is given by

$$x = ax_2m_2 + bx_1m_1 = 5 \times 8 \times 1 + 13 \times 7 \times 7 = 677$$

## 17.2 Complex Numbers

**Definition 17.1** (9.1.3). A **complex number** is an expression of the form a + bi where a and b are real numbers. The real number a is called the *real part* of a + bi, denoted as  $\Re(a + bi)$ . And the real number b is called the *imaginary part* of a + bi, denoted as  $\Im(a + bi)$ .

#### **Definition 17.2.**

$$i^2 = -1$$

Remark 17.2. Shorthands for complex numbers

- $\bullet \ a + i0 = a$
- 0 + ib = ib
- 0 + i0 = 0

#### **Remark 17.3.**

$$\mathbb{C} \subset \mathbb{R}$$

**Definition 17.3.** Arithmetic on complex numbers are defined as following,

• Addition(+ :  $\mathbb{C}^2 \to \mathbb{C}$ ) is defined as

$$(a+ib) + (c+id) := (a+c) + i(b+d)$$

• Multiplication( $\times : \mathbb{C}^2 \to \mathbb{C}$ ) is defined as

$$(a+ib) \times (c+id) := (ac-bd) + i(ad+bc)$$

**Proposition 17.1.** Let  $z = a + ib \in \mathbb{C}$  be a complex number, the **multiplication inverse** of z is given by

$$\frac{1}{a+ib} = \frac{a-ib}{a^2+b^2}$$

Proof.

$$\frac{1}{a+ib} = \frac{1}{a+ib} \times \frac{a-ib}{a-ib} = \frac{a-ib}{(a^2+b^2)+i(ab-ab)} = \frac{a-ib}{a^2+b^2}$$

# 18 Lecture 18 Oct. 19 2018 Complex Numbers

**Definition 18.1.** The **division** on complex numbers is equivalent to multiplying a complex number by its inverse, and is defined as

$$\frac{c+id}{a+id} = (c+id)\frac{a-ib}{a^2+b^2}$$

Notation 18.1. The set of all complex numbers is denoted as

$$\mathbb{C} = \{a + ib : a, b \in \mathbb{R}, i^2 = -1\}$$

**Remark 18.1.** Notice that  $\mathbb{C}$  is closed under the 4 basic operations of arithmetics. Anything like this is called a **field**.

**Example 18.1.** The irrational set  $\mathbb{Q}^c$  is *not* a field.

**Definition 18.2.** The **complex conjugate** of a complex number a + ib is defined as

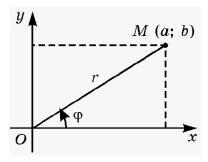
$$a - ib$$

**Definition 18.3.** The **modulus** of a complex number a + ib is defined as

$$|a+ib|=\sqrt{a^2+b^2}\in\mathbb{R}_{\geq 0}$$

# 18.1 The Geometric Representation of Complex Numbers

Any complex could be represented as a vector in a 2-dimensional coordinate, with real line on the x axis and imaginary line on the y axis.



**Remark 18.2** (Geometrical Interpretation). The **modulus** is the *distance* from the point to the origin.

**Remark 18.3** (Geometrical Interpretation). The **conjugate** is the *reflection* of the point about the real (x) axis.

#### 18.2 Polar Coordinates

#### 18.2.1 Coordinate Conversion

Consider a complex number represented by (a, b) in Cartesian coordinate and  $(r, \theta)$  in polar coordinate.

Cartesian	Polar
(a,b)	$(\sqrt{a^2+b^2}, \arctan(\frac{b}{a}))$
$(r\cos(\theta), r\sin(\theta))$	$(r,\theta)$

**Remark 18.4.**  $\arctan(\frac{b}{a})$  gives multiple solutions due to the periodicity of tan. We need to use the signs of real and imaginary values to determine which value of  $\arctan \frac{b}{a}$  to take.

**Example 18.2.** Consider 1 + i, it could be represented as (1, 1) in Cartesian coordinate. Converting it into polar coordinates gives  $(\sqrt{2}, \frac{\pi}{4})$ . Converting back gives

$$\begin{aligned} 1+i &= \sqrt{2}(\cos(\frac{\pi}{4})+i\sin(\frac{\pi}{4})) \\ &= \sqrt{2}(\cos(\frac{\pi}{4}+2\pi k)+i\sin(\frac{\pi}{4}+2\pi k)), \ \forall k \in \mathbb{Z} \end{aligned}$$

#### **18.2.2** Multiplication in Polar Coordinates

Consider the product of two complex numbers  $r_1(\cos(\theta_1)+i\sin(\theta_1))$  and  $r_2(\cos(\theta_2)+i\sin(\theta_2))$ :

$$r_1(\cos(\theta_1) + i\sin(\theta_1)) \times r_2(\cos(\theta_2) + i\sin(\theta_2))$$

$$= r_1 r_2 \Big[ \Big( \cos(\theta_1) \cos(\theta_2) - \sin(\theta_1) \sin(\theta_2) \Big) + i \Big( \cos(\theta_1) \sin(\theta_2) + \sin(\theta_1) \cos(\theta_2) \Big) \Big]$$
By triangle inequality
$$= r_1 r_2 \Big( \cos(\theta_1 + \theta_2) + i\sin(\theta_1 + \theta_2) \Big)$$

**Example 18.3.** Every complex number has a square root.

Proof. Let  $z = r(\cos(\theta) + i\sin(\theta)) \in \mathbb{C}$ . Consider  $w = \sqrt{r}(\cos(\frac{\theta}{2}) + i\sin(\frac{\theta}{2}))$ By above result, we could easily verify that  $w^2 = z$ . Notice that,  $w = \sqrt{r}(\cos(\frac{\theta}{2} + \pi) + i\sin(\frac{\theta}{2} + \pi))$  is also a square root of z.

#### 19 Lecture 19 Oct. 22 2018

#### 19.1 De Moivre's Theorem

**Theorem 19.1.** (De Moivre's Theorem) Let  $z = r[\cos(\theta) + i\sin(\theta)] \in \mathbb{C}$ , and the  $n^{th}$  power of z is given by

$$(r[\cos(\theta) + i\sin(\theta)])^n = r^n[\cos(n\theta) + i\sin(n\theta)], \ \forall n \in \mathbb{N}$$

*Proof.* (By induction)

**Base Case** for n = 1, obviously  $z^1 = z$ 

**Inductive Step** let  $k \in \mathbb{N}$ ,

suppose  $z^k = r^k [\cos(k\theta) + i\sin(k\theta)]$ 

Consider  $z^{k+1}$ ,

$$r^{k+1} = r^k [\cos(k\theta) + i\sin(k\theta)] \times r[\cos(\theta) + i\sin(\theta)]$$

$$= r^{k+1} \Big[ (\cos(k\theta)\cos(\theta) - \sin(k\theta)\sin(\theta)) + i(\cos(k\theta)\sin(\theta) + \sin(k\theta)\cos(\theta)) \Big]$$
By Triangle Identity
$$= r^{k+1} \Big[ \cos((k+1)\theta) + i\sin((k+1)\theta) \Big]$$

We could then conclude what the theorem stated by principle of mathematical induction.

**Example 19.1.** Calculate  $(1+i)^8$ .

Solution. 1 + i can be written as (1, 1) in Cartesian coordinate.

Then it can be converted into Polar coordinate as

$$\sqrt{2}\Big(\cos(\frac{\pi}{4}) + i\sin(\frac{\pi}{4})\Big)$$

Then by De Moivre's theorem,

$$\left(\sqrt{2}\left(\cos(\frac{\pi}{4}) + i\sin(\frac{\pi}{4})\right)\right)^{8}$$

$$= (\sqrt{2})^{8}\left(\cos(\frac{\pi}{4} \times 8) + i\sin(\frac{\pi}{4} \times 8)\right)$$

$$= 16 \times \left(\cos(2\pi) + i\sin(2\pi)\right)$$

$$= 16(\cos(0) + i\sin(0))$$

$$= 16$$

Geometrically Interpretation: rotates the vector anti-clockwise by  $(n-1)\theta$  and enlarge the magnitude by factor of n.

## 19.2 Roots of Unity

**Example 19.2.** Find all roots of  $z^2 = 1$ , where  $z \in \mathbb{C}$ .

Solution. In polar coordinates, let  $z = r(\cos \theta + i \sin \theta)$ . Thus by De Moivre's Theorem,  $z^2 = r^2(\cos(2\theta) + i \sin(2\theta))$ . And  $1 = 1(\cos(0 + 2k\pi) + \sin(0 + 2k\pi))$ ,  $k \in \mathbb{Z}$  in polar coordinate. Solving the equation  $z^2 = 1$  gives

$$\begin{cases} r^2 = 1 \\ 2\theta = 2k\pi, \ k \in \mathbb{Z} \end{cases}$$

We can conclude that r = 1 since it represents a *distance* and  $r \in \mathbb{R}_{\geq 0}$ .

- k = 0: r = 1,  $\theta = 0 \rightarrow 1(\cos(0) + i\sin(0)) = 1$
- k = 1: r = 1,  $\theta = \pi \to 1(\cos(\pi) + i\sin(\pi)) = -1$
- k = 2: r = 1,  $\theta = 2\pi \rightarrow 1(\cos(2\pi) + i\sin(2\pi)) = 1$

From the repeating pattern we can conclude that  $\forall k \in \mathbb{Z}^5$ 

$$z = 1(\cos(\pi k) + i\sin(\pi k)) = \pm 1$$

**Example 19.3.** Find all roots of  $z^n = 1$ 

Solution.

$$z^{n} = r^{n} [\cos(n\theta) + i \sin(n\theta)]$$

$$1 = 1(\cos(2k\pi) + i \sin(2k\pi))$$

$$\implies r = 1$$

$$\land n\theta = 2k\pi \iff \theta = k\frac{2\pi}{n}$$

Consider cases

- k = 0:  $r = 1, \theta = 0$
- k = 1:  $r = 1, \ \theta = \frac{2\pi}{n}$

<sup>&</sup>lt;sup>5</sup>The case k < 0 is covered by symmetry.

• k = 2: r = 1,  $\theta = 2\frac{2\pi}{n}$ 

• k = 3: r = 1,  $\theta = 3\frac{2\pi}{n}$ 

Until k = n, we have  $r = 1 \land \theta = n \frac{2\pi}{n} = 2\pi$ , where  $z|_{k=n} = z|_{k=0}$  and the root starts repeating.

There are n roots in total,

$$z = \cos(k\frac{2\pi}{n}) + \sin(k\frac{2\pi}{n}), \ k \in \{0, 1, \dots, n-1\}$$

**Example 19.4.** Solve  $z^3 = 1$ 

**Example 19.5.** Solve  $z^4 = 1$ 

**Geometrically Interpretation** Divides the unit ball into n equal slices.

**Example 19.6.** Solve  $z^3 = 2 + 2i$ 

Solution. In polar coordinate,

$$2 + 2i = \sqrt{8} \left( \cos(\frac{\pi}{4} + 2k\pi) + i \sin(\frac{\pi}{4} + 2k\pi) \right)$$

We have to solve

$$3\theta = \frac{\pi}{4} + 2k\pi, \ k \in \mathbb{Z}$$

$$\implies \theta = \frac{\pi}{12} + k\frac{2\pi}{3}, \ k \in \mathbb{Z}$$

And clearly  $r = \sqrt{2}$ . And roots are fond by plugging in k with 0, 1, 2.

$$\begin{cases} z_1 = \sqrt{2}(\cos(\frac{\pi}{12}) + i\sin(\frac{\pi}{12})) \\ z_2 = \sqrt{2}(\cos(\frac{\pi}{12} + \frac{2\pi}{3}) + i\sin(\frac{\pi}{12} + \frac{2\pi}{3})) \\ z_3 = \sqrt{2}(\cos(\frac{\pi}{12} + \frac{4\pi}{3}) + i\sin(\frac{\pi}{12} + \frac{4\pi}{3})) \end{cases}$$

#### 20 Lecture 20. Oct 24 2018

**Theorem 20.1** (The Fundamental Theorem of Algebra). Every *non-constant* polynomial (with complex coefficients) has a complex root. i.e. for

$$p(z) = a_n z^n + a_{n-1} z^{n-1} + \dots + a_1 z + a_0, \quad a_i \in \mathbb{C}, \ n \ge 1$$

there exists  $r \in \mathbb{C}$  such that p(r) = 0.

**Example 20.1.** Let  $p(x) = x^3 - 3x^2 - 9x + 27$ ,  $p(x) = (x - 3)^2(x + 3)$ .

**Interpretation** *Linear polynomials (polynomials of degree 1) are building blocks of all polynomials.* 

**Theorem 20.2** (Long Division). Suppose p(z) is a non-constant polynomial and  $r \in \mathbb{C}$ , there exists a polynomial  $q(z) \in \mathcal{P}$  and  $c \in \mathbb{C}$  such that

$$p(z) = q(z)(z - r) + c$$

where c is the *reminder* of long division.

**Definition 20.1.** A polynomial f(z) is a **factor** of another polynomial p(z) if

$$\exists q(z) \in \mathcal{P}, \ s.t. \ p(z) = f(z)q(z)$$

**Theorem 20.3** (Factor Theorem). A complex number r is a root of a polynomial p(z) if and only if (z - r) is a factor of p(z).

*Proof.* ( $\iff$ ) Suppose (z - r) is a factor.

By definition of factor,  $\exists q(z) \in \mathcal{P}$  such that p(z) = q(z)(z - r).

Plugging in r gives p(r) = q(r)(r - r) = 0 and suggests r is a root of p(z).

 $(\Longrightarrow)$  Suppose r is a root of p(z).

By the theorem of long division,  $\exists q(z) \in \mathcal{P}$  and  $c \in \mathbb{C}$  satisfying

$$p(z) = q(z)(z - r) + c.$$

Plugging in z = r gives p(r) = q(r)(r - r) + c = 0, which implies c = 0.

That's 
$$p(z) = q(z)(z - r)$$
.

**Theorem 20.4** (Extended Fundamental Theorem of Algebra). A non-zero<sup>6</sup> polynomial of degree n has exactly n roots, counting multiplicities.

*Proof.* Let  $p(z) \in \mathcal{P}$  with degree  $n \ge 0$  and suppose p(z) is non-zero.

Case 1: n = 0, then p(z) has 0 roots.

Case 2:  $n \ge 1$ , by the fundamental theorem of algebra,

<sup>&</sup>lt;sup>6</sup>For zero polynomial, it has infinitely many roots.

p(z) has a root  $r_1 \in \mathbb{C}$ .

By factor theorem,  $\exists q_1(z) \in \mathcal{P}(\mathbb{C}), \ s.t. \ p(z) = (z - r_1)q_1(z).$ 

Note that  $q_1(z)$  has degree of n-1.

If  $n - 1 \ge 1$ , repeating above argument and we have

 $\exists r_2 \in \mathbb{C}, \ \exists q_2(z) \in \mathcal{P}(\mathbb{C}), \ s.t. \ q_1(z) = (z - r_2)q_2(z).$ 

Note that  $q_2(z)$  has degree of n-2.

Equivalently  $p(z) = (z - r_1)(z - r_2)q_2(z)$ .

Iterating till  $q_i(z)$  has degree 0 (i.e. constant), this will be achieved after exactly n iterations.

Aggregately, we can factorize p(z) into

$$p(z) = (z - r_1)(z - r_2) \dots (z - r_n)q_n$$

where  $q_n$  is a constant.

Obviously there are n (possibly repeating) roots, namely  $r_1, r_2, \ldots, r_n$ .

# 21 Lecture 21. Oct 29 2018

Lemma 21.1 (Triangle Inequality).

$$|z_1 + z_2| \le |z_1| + |z_2|, \ \forall z_1, z_2 \in \mathbb{C}$$

Lemma 21.2 (Extended version of triangle inequality).

$$|\sum_{i=1}^n z_i| \le \sum_{i=1}^n |z_i|, \ \forall (z_i) \in \mathbb{C}^n$$

**Definition 21.1.** A **closed curve in the plane** is a continuous function mapping from  $[0, 2\pi]$  to  $\mathbb{C}$  such that its values at 0 and  $2\pi$  are the same.

**Definition 21.2.** If  $\phi(t) : [0, 2\pi] \to \mathbb{C}$  is a closed curve that does not go through the origin, its **winding number** is the number of times a vector from the origin to a point on the curve winds around the origin as t goes from 0 to  $2\pi$ .

Example 21.1. Consider

$$\phi(t) = f(t) + i(g)$$

where  $f, g : [0, 2\pi] \to \mathbb{R}$  are continuous. Then  $\phi(t)$  is continuous.

Example 21.2. Consider

$$\phi(t) = \cos(t) + i\sin(t)$$

the function above is a closed curve with winding number +1.

**Remark 21.1.** If points on the curve go around the origin *anti-clockwise* as t goes from 0 to  $2\pi$ , then we consider the winding number to be *negative*.

**Example 21.3.** Curve  $\phi(t) = \cos(3t) + i\sin(3t)$  has winding number +3.

**Example 21.4.** Curve  $\phi(t) = 27\cos(4t) + 27i\sin(4t)$  has winding number +4.

**Example 21.5.** Curve  $\phi(t) = \sin(t) + i\cos(t)$  has winding number -1.

**Example 21.6.** A non-zero constant (e.g.  $\phi(t) = 3 + 4i$ ) is closed and not passing the origin, it has winding number 0.

Remark 21.2. The notation of winding number only apply to closed curves that do **not** passes the origin.

### 21.1 Proof of the Fundamental Theorem of Algebra

*Proof.* Idea: prove by contradiction.

Suppose p(z) is a non-constant polynomial with no roots. i.e.

$$p(z) \neq 0, \ \forall z \in \mathbb{C}$$

and degree of p(z) = n > 0.

For each radius R > 0 define

$$\phi_R(t) := R(\cos(t) + i\sin(t))$$

Then for each R > 0, let

$$p_R(t) := p(\phi_R(t))$$

note that  $p_R(t): [0, 2\pi] \to \mathbb{C}$  and it's a closed curve.

Also note that since  $p_R(t) \neq 0 \ \forall t \in [0, 2\pi], \ p_R(t)$  does not go through the origin. We will show that

- 1. If R is large enough then the winding number of  $p_R(t) = deg(p(z))$ .
- 2. If *R* is small enough then the winding number of  $p_R(t) = 0$ .

But the winding number of  $p_R(t)$  is a continuous function of R and it has co-domain of integers. Then it must be the case that  $p_R(t)$  is constant, but this contradicts our assumption that deg(p(z)) > 0.

# 22 Lecture 22. Oct 31 2018

**Recall** the outline of proving the Fundamental Theorem of Algebra. Suppose p(z) is a non-=constant polynomial with no roots. i.e.

$$p(z) \neq 0 \quad \forall z \in \mathbb{C}$$

Let

$$p_R(t) := p(\phi_R(t))$$

where  $\phi(t) : [0, 2\pi] \to \mathbb{C} = R(\cos(t) + i\sin(t)).$ 

And we well show

- 1. R is large  $\implies$  winding number of  $p_R(t) = degree(p(z))$
- 2. *R* is small  $\implies$  winding number of  $p_R(t) = 0$ .

*Proof.* Let  $q(z) = z^n$ , where n is the degree of polynomial p(z). Let  $L_R(t) = q(\phi_R(t))$ .

Note that

$$L_R(t) = q(\phi_R(t))$$

$$= q(R(\cos(t) + i\sin(t)))$$

$$= R^n(\cos(nt) + i\sin(nt))$$

so  $L_R(t)$  has winding number n.

**Lemma 22.1.** Let L(t) and M(t) be 2 closed curves not passing through the origin. Suppose

$$|L(t) - M(t)| < |L(t)| \quad \forall t \in [0, 2\pi]$$

then L(t) and M(t) have the same winding number.

We are **not** going to prove this lemma.

*Proof. Proposition 1.* Since  $L_R(t) = \phi_R(t)^n$ ,

Suppose

$$p(\phi_R(t)) = a_n z^n + a_{n-1} z^{z-1} + \dots + a_1 z + a_0$$

WLOG, assume  $a_n = 1$ .

Then

$$p_R(t) = \phi_R(t)^n + a_{n-1}\phi_R(t)^{n-1} + \dots + a_1\phi_R(t) + a_0$$

and

$$|L_{R}(t) - p_{R}(t)| = |a_{n-1}\phi_{R}(t)^{n-1} + \dots + a_{1}\phi_{R}(t) + a_{0}|$$

$$\leq |a_{n-1}\phi_{R}(t)^{n-1}| + \dots + |a_{0}|$$

$$= |a_{n-1}||\phi_{R}(t)|^{n-1} + |a_{n-2}||\phi_{R}(t)|^{n-2} + \dots + |a_{1}||\phi_{R}(t)| + |a_{0}|$$

$$= |a_{n-1}|R^{n-1} + |a_{n-2}|R^{n-2} + \dots + |a_{1}|R + |a_{0}|$$

$$\text{Choosing } R > \max\{1, \sum_{i=1}^{n-1} |a_{i}|\}$$

$$< |a_{n-1}|R^{n-1} + |a_{n-2}|R^{n-1} + \dots + |a_{1}|R^{n-1} + |a_{0}|R^{n-1}|$$

$$= R^{n-1} \sum_{i=1}^{n-1} |a_{i}|$$

$$< R^{n} = |L_{R}(t)|$$

Thus we have shown that

$$|L_R(t) - p_R(t)| < |L_R(t)|, \quad \forall t \in [0, 2\pi]$$

by choosing R large enough. By previous lemma, we conclude that  $p_R(t)$  has the same winding number as  $L_R(t)$ , which is n.

*Proof. Proposition 2.* Note  $p(0) = a_0 \neq 0$  since we assumed p has no roots. Since p(z) is a polynomial so its continuous. (of course near 0)

$$\forall \epsilon > 0, \ \exists \delta > 0 \ s.t. \ 0 < |z - 0| < \delta \implies |p(z) - p(0)| < \epsilon$$

Since  $p(0) \neq 0$  and the quadrant(excluding axes) containing  $a_0$  is open.

There exists  $\epsilon$  such that all points z in  $\mathcal{B}(\epsilon, a_0)$  are in that quadrant.

There exists  $\delta > 0$  satisfying the continuity definition above and we choose

$$R = \frac{\delta}{2}$$

Then all z in set  $\{\phi_R(t): t \in [0, 2\pi]\}$  are mapped into  $\epsilon$ ,  $\dashv$ , and of course in the quadrant containing  $a_0$ .

Therefore the winding number of  $p_R(t)$  is 0.

altogether with the fact that winding number of  $p_R(t)$  is a continuous function from  $R_{>0}$  to integers, we conclude that  $p_R(t)$  is constant.

This conclusion contradicts our assumption that p(z) is non-constant, i.e.  $n \neq 0$ . Thus p(z) has root.

#### 23 Lecture 23. Nov 2 2018

**Proposition 23.1.** The winding number transformation of  $p_R(t) = p(\phi_R(t))$ ,  $W : \mathbb{R}_{>0} \to \mathbb{Z}$  is continuous in R.

*Proof.* For small enough  $\epsilon > 0$ ,

Consider

$$|p_{R+\epsilon}(t) - p_R(t)| < |p_R(t)|, \ \forall t \tag{1}$$

we will show (1) is true for sufficiently small  $\epsilon > 0$ .

By lemma 22.1,  $p_{R+\epsilon}(t)$  and  $p_R(t)$  have the same winding number.

Let

$$p(z) = z^{n} + a_{n-1}z^{n-1} + \dots + a_{1}z + a_{0}$$
(2)

Then

$$|p_{R+\epsilon}(t) - p_R(t)| = |(\phi_{r+\epsilon}(t)^n + a_{n-1}\phi_{r+\epsilon}(t)^{n-1} + \dots) - (\phi_R(t)^n + \dots)|$$
 (3)

$$= \left| (\phi_{r+\epsilon}(t)^n - \phi_R(t)^n) + a_{n-1}(\phi_{r+\epsilon}(t)^{n-1} - \phi_R(t)^{n-1}) + \dots \right| \tag{4}$$

$$= \left| [(R + \epsilon)^n - R^n] e^{int} + a_{n-1} [(R + \epsilon)^{n-1} - R^{n-1}] e^{i(n-1)t} + \dots \right|$$
 (5)

$$\leq \left| (R + \epsilon)^{n} - R^{n} \right| \left| e^{int} \right| + \left| a_{n-1} \right| \left| (R + \epsilon)^{n-1} - R^{n-1} \right| \left| e^{i(n-1)t} \right| + \dots + \left| a_{1} \right| \left| e^{it} \right| \tag{6}$$

note that

$$|e^{ijt}| = 1, \quad \forall j \in \{1, 2, \dots, n\}$$
 (7)

Thus

$$\left| p_{R+\epsilon}(t) - p_R(t) \right| \le \sum_{j=1}^n \left| (R+\epsilon)^j - R^j \right| \tag{8}$$

Note that we can make  $|(R + \epsilon)^k - R^k|$  as small as we want by specifying a sufficiently small  $\epsilon$  since  $x^k$  is continuous.

**Definition 23.1.** A set *S* is **finite** if there exists some  $n \in \mathbb{N}$  such that the elements of *S* can be paired with the elements in set  $\{1, 2, ..., n\}$ . Equivalently, we can label the elements of *S* as  $s_1, s_2, ..., s_n$ .

**Definition 23.2.** A set is **infinite** if it is not finite.

**Definition 23.3.** Two sets S and  $T^7$  have the same **cardinality** if and only if there exists a *bijection* between them. Written as |S| = |T|.

 $<sup>^{7}</sup>S$  and T are **not** necessarily finite.

#### 24 Lecture 24. Nov 12 2018

**Example 24.1** (Infinite Sets with Same Cardinality). Let  $S = \mathbb{N}$  and  $T = \{2, 4, 6, ...\}$ , easy to construct a bijective mapping  $f : S \to T$  defined as f(n) = 2n to show S and T have the same cardinality.

**Example 24.2** (Infinite Sets with Same Cardinality). Let  $S = \mathbb{N}$  and  $T = \{2, 3, 4, ...\}$ , easy to construct a bijective mapping  $f : S \to T$  defined as f(n) = n + 1 to show S and T have the same cardinality.

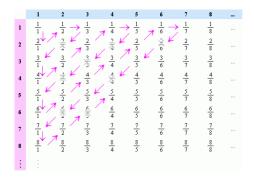
**Remark 24.1** (How to Prove Same Cardinality with Natural Numbers). If  $|\mathbb{N}| = |T|$ , then  $\exists$  bijection  $f : \mathbb{N} \to T$ . That's we can *enumerate* elements in T with f(n):

$$t_1 = f(1), t_2 = f(2), \dots \forall t_i \in T$$

**Definition 24.1.** Let T be a set, if  $|T| = |\mathbb{N}|$ , then T is **countably infinite** and written as  $|T| = \aleph_0$ .

**Definition 24.2.** A set is **countable** if it is either finite or countably infinite.

**Theorem 24.1.** Set  $S = \mathbb{Q}^+$  is countable infinite.



Proof.

*Proof.* (alternative.) Note that the Cartesian product of countably infinite sets is countably infinite.

Then, for all  $\frac{p}{q} \in \mathbb{Q}^+$ , where  $p, q \in \mathbb{N}$  and p, q are relatively prime.

Setup bijection  $\phi : \mathbb{Q}^+ \to \mathbb{N}^2$ , defined as  $\phi(\frac{p}{q}) = (p, q)$ .

So  $|\mathbb{Q}^+| = |\mathbb{N}^2|$ .

Thus  $\mathbb{Q}^+$  is countably infinite.

**Theorem 24.2.** Let [0, 1] be the set of all real numbers between 0 and 1. [0, 1] is not infinitely countable.

*Proof.* (Prove by contradiction)

Suppose we can list all real numbers between 0 and 1 as

$$a_1 = 0.a_{11}a_{12}a_{13} \dots$$
  
 $a_2 = 0.a_{21}a_{22}a_{23} \dots$   
 $a_3 = 0.a_{31}a_{32}a_{33} \dots$   
:

Consider another real number in [0, 1] constructed as following

$$x = 0.x_1x_2x_3...$$
 where  $x_i = \begin{cases} 6 \text{ if } a_{ii} = 5\\ 5 \text{ otherwise} \end{cases}$ 

Note that the first decimal of  $x(x_1)$  is not the same as the first decimal of  $a_1(a_{11})$ , so  $x \neq a_1$ .

Similarly, the second decimal of  $x(x_2)$  is not the same as the second decimal of  $a_2(a_{22})$ , so  $x \neq a_2$ ,

It is easy to show that  $x \neq a_i$ ,  $\forall i$ . So x is a real number between 0 and 1 not included in the table above.

Contradicting the assumption that we could list all real numbers between 0 and 1 in a table.

Thus [0, 1] is not infinitely countable.

#### 25 Lecture 25. Nov 14 2018

**Notation 25.1** (the Cardinality of continuum). |[0,1]| = C

**Definition 25.1.**  $|S| \le |T|$  if there exists a subset  $T_0 \subseteq T$  such that  $|T_0| = |S|$ . Or, equivalently, there exists an injection maps from S to T.

**Definition 25.2.**  $|S| < |T| \text{ if } |S| \le |T| \text{ and } |S| \ne |T|.$ 

**Proposition 25.1.**  $|\mathbb{N}| < |[0, 1]|$  (i.e.  $\aleph_0 < C$ )

*Proof.* We've already shown that  $|\mathbb{N}| \neq |[0, 1]|$ .

Consider injection  $f: \mathbb{N} \to [0, 1]$  defined as  $f(n) = \frac{1}{n}$ .

Therefore  $|\mathbb{N}| \leq |[0, 1]|$ .

Or equivalently the subset of [0,1] defined as  $\{\frac{1}{n}: n \in \mathbb{N}\}$  has the same cardinality as  $\mathbb{N}$ .

Thus, by definition,  $|\mathbb{N}| < |[0, 1]|$ .

**Theorem 25.1** (Schödre-Bernstein-Cantor Theorem). Let *S* and *T* be two sets then

$$|S| \le |T| \land |S| \ge |T| \implies |S| = |T|$$

Proof.

**Proposition 25.2.** Let  $a, b, c, d \in \mathbb{R}$  satisfying a < b and c < d, then

$$|[a,b]| = |[c,d]| = C$$

Every closed interval has the same cardinality.

*Proof.* Consider mapping  $f(x) = (d - c)\frac{x-a}{b-a} + c$ .

Obviously,  $f : [a, b] \rightarrow [c, d]$  and bijective.

And therefore it's inverse is a bijection from [c, d] to [a, b].

Thus those two closed intervals have the same cardinality.

**Proposition 25.3.**  $|\mathbb{R}| = |(-\frac{\pi}{2}, \frac{\pi}{2})|$ .

*Proof.* Consider bijection  $f(x) := \tan(x)$ .

**Proposition 25.4.** |[0, 1]| = |(0, 1)|.

*Proof.* Step 1. Consider bijection f(x) := x and obviously  $|(0, 1)| \le |[0, 1]|$ .

Step 2. As shown before, all closed interval have the same cardinality. Thus  $|[0,1]| = |[\frac{1}{4}, \frac{1}{2}]|$ . And clearly  $[\frac{1}{4}, \frac{1}{2}] \subseteq (0,1)$ . So  $|[0,1]| = |[\frac{1}{4}, \frac{1}{2}]| \le |(0,1)|$ .

By Schödre-Bernstein-Cantor theorem, |[0, 1]| = |(0, 1)|

**Proposition 25.5.** Above result can be generalized to arbitrary open and closed intervals, i.e.

$$|[a,b]| = |(c,d)|$$