

# CSC165 Lecture notes

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## Info.

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## 1 Lecture 2 Jan.15 2018

### 1.1 Predicate Logic

*Allow you to have a domain of objects that you want to talk about, we want to be express reason about this domain.*

**Predicate** Simplest kind of Predicate Logical Formula is called a **predicate**. A predicate is a function with range  $\{0, 1\}$ .

#### Examples

1. *Less-than-or-equal-to:*  $\leq: \mathbb{Z} \times \mathbb{Z} \rightarrow \{0, 1\}$
2. *Equality:*  $=: \mathbb{Z} \times \mathbb{Z} \rightarrow \{0, 1\}$
3. *Prime:*  $Prime: \mathbb{N} \rightarrow \{0, 1\}$
4. Define:  $R: \{a, b\} \times \{1, 2, 3\} \rightarrow \{0, 1\}$  as  $R(a, 1) = R(b, 1) = R(c, 1) = 1, R = 0$  otherwise.

When you specify/define a predicate you have to specify the domain.

**Quantifiers** Introduce two quantifiers, **exists**:  $\exists$  and **for all**:  $\forall$ , that let us express,

$\exists x \in \text{DOMAIN}$ : There is at least one element in domain of predicate that is true.

equivalently, represent as  $\vee$ ,

$$"\exists" \equiv p(x_0) \vee p(x_1) \vee p(x_2) \dots$$

$\forall x \in \text{DOMAIN}$ : All element in domain of predicate satisfy the predicate.

equivalently, represented as  $\wedge$ ,

$$"\forall" \equiv p(x_0) \wedge p(x_1) \wedge p(x_2) \dots$$

### Negation of quantifier statements

$$\neg(\exists x \in \mathbb{D}, \text{ s.t. } p(x)) \equiv \forall x \in \mathbb{D}, \neg p(x)$$

$$\neg(\forall x \in \mathbb{D}, \text{ s.t. } p(x)) \equiv \exists x \in \mathbb{D}, \neg p(x)$$

**Nested quantifier** more than one variable quantified.

#### Example

*For every natural number  $x$ , if  $x$  is a power of 2 then  $2x$  is a power of 2.*

$$\forall x \in \mathbb{N}, (\exists k \in \mathbb{N} \text{ s.t. } x = 2^k) \implies (\exists k' \in \mathbb{N} \text{ s.t. } 2^{k'} = 2x)$$

## 2 Lecture 3 Jan.17 2018

**Another example** There are infinitely many natural number that are even.

$$\text{Even}(\dots): \exists y \in \mathbb{N}, x = 2y$$

$$\forall x \in \mathbb{N}, \exists y \text{ s.t. } y > x \wedge \text{Even}(y)$$

$$Q(x) : \text{some predicate}$$

$$Q(x) : \mathbb{N} \rightarrow \{0, 1\}$$

$$\forall x \in \mathbb{N}[Q(x) \implies \exists y \in \mathbb{N}[y > x \wedge Q(y)]]$$

Does this express *there are infinitely many numbers that satisfy  $Q$* ?

Not, consider a  $Q$  that is false for all  $x$ , the statement is vacuous truth, but it does not express what we want.

Fix it.

$$\exists z \in \mathbb{N}[Q(z)] \wedge \forall x \in \mathbb{N} \exists y \in \mathbb{N}[y > x \wedge Q(y)]$$

To ensure that there are some elements satisfy  $Q$

**Tips**

1. Make sure you write down a predicate logic formula with **correct syntax**.
2. Use different variable names for different quantities.
3. When you quantify ( $\exists x$  or  $\forall y$ ) you must specify the **set**. e.g. We want to quantify over all  $x \in \mathbb{N}$  such that  $x \geq 5$ .

(a) Method 1: Predefine your set.

$$\text{Let } S = \{x \mid x \in \mathbb{N}, x \geq 5\} \quad \forall x \in S, x \geq 3$$

(b) Method 2: By implication.

$$\forall x \in \mathbb{N}, \quad x \geq 5 \implies x \geq 3$$

**Note** We will spend a lot of time defining new predicates using predicate logic and then reasoning about them.

**Example: Divisibility** Let  $n, d \in \mathbb{Z}$ , we say that  $d$  **divides**  $n$  or  $n$  is **divisible** by  $d$  iff

$$\exists k \in \mathbb{Z} \text{ s.t. } n = d * k$$

$$\text{Divides}(d, n) : \exists k \in \mathbb{Z} [n = d * k]$$

this formula is a predicate logic, some of variables ( $k$ ) are quantified and others ( $n, d$ ) are not quantified, called **free variable**. This formula since it has free variables, it represents a predicate.

A formula with no free variable is called a **sentence**, sentences are true or false.

**Let's express** For every integer  $x$ , if  $x$  divides 10 then it also divide 100.

$$\forall x \in \mathbb{Z} [\text{Divides}(x, 10) \implies \text{Divides}(x, 100)]$$

equivalently,

$$\forall x \in \mathbb{Z} [\exists k \in \mathbb{Z} [10 = k * x] \implies \exists k' \in \mathbb{Z} [100 = k' * x]]$$

**Note** Proposition is a special type of predicate but it's **not** a sentence.

## 3 Lecture 4 Jan.22 2018

### 3.1 Introduction to proofs

**Proof** A **proof** is a logical argument that convinces another person that a statement is true. Can also have a **disproof** showing that a statement is false.

1. Write down what we want to prove using language of first order logic.
2. Introducing variable(s).
3. Write body of proof.

**Example** Prove that every natural number  $n$  satisfies the inequality  $n^2 + 3n + 7 \geq 4$ ,

**Statement in first order logic:**

$$\forall n \in \mathbb{N}, n^2 + 3n + 7 \geq 4$$

*proof:*

**Introducing variable(s):**

Let  $n \in \mathbb{N}$

**Body of proof:**

Since  $n \in \mathbb{N}$ ,  $n \geq 0$

Therefore  $n^2 \geq 0$

Similarly, since  $n \in \mathbb{N}$ ,  $n \geq 0$ ,  $3n \geq 0$

$$\therefore n^2 + 3n + 7 \geq 0 + 0 + 4$$

■

**Example** Prove that for every natural number  $n$  greater than 20,  $n$  satisfies  $1.5n - 4 \geq 3$ .

$$\forall n \in \mathbb{N}, [ n > 20 \implies 1.5n - 4 > 3 ]$$

*proof.*

Let  $n \in \mathbb{N}$ , assume that  $n > 20$

Since  $n > 20$ ,  $1.5n > 1.5(20) = 30$

So  $1.5n > 30$

$$\therefore 1.5n - 4 > 26$$

$$\therefore 1.5n - 4 > 3$$

■

**(More complex)Example** Define a natural number to be a **Prime** numbers:

$$Divides(x, n) : \exists xk = n$$

$$Prime(n) : (n > 1) \wedge (\forall x \in \mathbb{N}, ((x \neq 1 \wedge x \neq n) \implies \neg Divides(x, n)))$$

Equivalently, take contrapositive

$$Prime(n) : (n > 1) \wedge (\forall x \in \mathbb{N}, Divides(x, n) \implies (x = 1 \vee x = n))$$

**Example** For every integer  $x$ ,  $x|x+1$  then  $x|5$

$$\forall x \in \mathbb{Z}, [ Divides(x, x+5) \implies Divides(x, 5) ]$$

$$\forall x \in \mathbb{Z}, [ (\exists k \in \mathbb{Z} \text{ s.t. } xk = x + 5) \implies (\exists k' \in \mathbb{Z} \text{ s.t. } xk' = x) ]$$

*proof.*

Let  $x \in \mathbb{Z}$

Assume  $k \in \mathbb{Z}$  is such that  $xk = x + 5$

Let  $k' = k - 1$

Then,  $k'x = (k - 1)x$

$$= kx - x$$

By assumption,

$$= 5$$

Therefore,  $\exists k' \in \mathbb{Z}, \text{ s.t. } k'x = 5$

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