

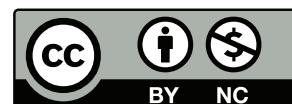
# ECO326 Advanced Microeconomic Theory

A Course in Game Theory

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**Readme** this note is based on the course content of *ECO326 Advanced Microeconomics - Game Theory*, this note contains all materials covered during lectures and mentioned in the course syllabus. However, notations, statements of theorems and proofs are following the book *A Course in Game Theory* by Osborne and Rubinstein, so they might be, to some extent, more mathematical than the required text for ECO326, *An Introduction to Game Theory*.

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# 1 Lecture 1. Games and Dominant Strategies

**Assumption 1.1** (pg.4). Assume that each decision-maker is *rational* in the sense that he is aware of his alternatives, forms expectation about any unknowns, has clear preferences, and chooses his action deliberately after some process of optimization.

**Definition 1.1** (pg.4). A model of **rational choice** consists

- A set  $A$  of *actions*.
- A set  $C$  of *consequences*.
- A *consequence function*  $g : A \rightarrow C$ .
- A *preference relation*  $\succsim$  on  $C$ .

**Definition 1.2** (pg.7). A **preference relation** is a complete reflexive and transitive binary relation.

**Definition 1.3** (11.1). A **strategic game** consists of

- a finite set of **players**  $N$ .
- for each player  $i \in N$ , an **actions**  $A_i \neq \emptyset$ .
- for each player  $i \in N$ , a **preference relation**  $\succsim_i$  defined on  $A \equiv \times_{i \in N} A_i$ .

and can be written as a triple  $\langle N, (A_i), (\succsim_i) \rangle$ .

**Definition 1.4** (pg.11). A strategic game  $\langle N, (A_i), (\succsim_i) \rangle$  is **finite** if

$$|A_i| < \aleph_0 \quad \forall i \in N$$

## 2 Lecture 2. Iterated Elimination and Rationalizability

### 2.1 Iterated Elimination of Strictly Dominated Strategies (Actions)

**Definition 2.1** (60.2). The set  $X \subseteq A$  of outcomes of a finite strategic game  $\langle N, (A_i), (u_i) \rangle$  **survives iterated elimination of strictly dominated actions** if  $X = \times_{j \in N} X_j$  and there is a collection  $((X_j^t)_{j \in N})_{t=0}^T$  of sets that satisfies the following conditions for each  $j \in N$ .

- $X_j^0 = A_j$  and  $X_j^T = X_j$ .
- $X_j^{t+1} \subseteq X_j^t$  for each  $t = 0, \dots, T-1$ .
- For each  $t = 0, \dots, T-1$  every action of player  $j$  in  $X_j^t \setminus X_j^{t+1}$  is strictly dominated in the game  $\langle N, (X_i^t), (u_i^t) \rangle$ , where  $u_i^t$  for each  $i \in N$  is the function  $u_i$  restricted to  $\times_{j \in N} X_j^t$ .
- No action in  $X_j^T$  is strictly dominated in game  $\langle N, (X_i^T), (u_i^T) \rangle$ .

**Proposition 2.1** (61.2). If  $X = \times_{j \in N} X_j$  survives iterated elimination of strictly dominated actions in a finite strategic game  $\langle N, (A_i), (u_i) \rangle$  then  $X_j$  is the set of player  $j$ 's rationalizable actions for each  $j \in N$ .

## 2.2 Rationalizability

**Definition 2.2** (pg. 54). A **belief** of player  $i$  (about the actions of the other players) is a probability measure,  $\mu_i$ , on  $A_{-i} = \times_{j \in N \setminus \{i\}} A_j$ .

**Definition 2.3** (59.1). An action of player  $i$  in a strategic game is a **never best response** if it is not a best response to any belief of player  $i$ .

**Definition 2.4** (59.2). The action  $a_i \in A_i$  of player  $i$  in the strategic game  $\langle N, (A_i), (u_i) \rangle$  is **strictly dominated** if there is a mixed strategy  $\alpha_i$  of player  $i$  such that

$$U_i(a_{-i}, \alpha_i) > u_i(a_{-i}, a_i)$$

for all  $a_{-i} \in A_{-i}$ , where  $U_i(a_{-i}, \alpha_i)$  is the payoff of player  $i$  if he uses the mixed strategy  $\alpha_i$  and the other players' vector of actions is  $a_{-i}$ .

## 3 Lecture 3. Nash Equilibrium

**Definition 3.1** (14.1). A **Nash equilibrium of a strategic game**  $\langle N, (A_i), (\succsim_i) \rangle$  is a profile  $a^* \in A$  of actions with property that for every player  $i \in N$

$$(a_i^*, a_{-i}^*) \succsim_i (a_i, a_{-i}^*) \quad \forall a_i \in A_i$$

**Definition 3.2** (pg.15). The **best-response function** for a player  $i$  is defined as

$$B_i(a_{-i}) = \{a_i \in A_i : (a_i, a_{-i}) \succsim_i (a'_i, a_{-i}) \quad \forall a'_i \in A_i\}$$

**Remark 3.1.** The best-response of  $a_{-i}$  can be written as

$$B_i(a_{-i}) = \bigcap_{a'_i \in A_i} \{a_i \in A_i : (a_i, a_{-i}) \succsim_i (a'_i, a_{-i})\}$$

where each of them is the upper contour set of  $a'_i$ .

Thus, if  $\succsim_i$  is quasi-concave, then  $B_i(a_{-i})$  is an intersection of convex sets and therefore itself convex.

**Remark 3.2** (pg.15). So a Nash equilibrium is a profile  $a^* \in A$  such that

$$a_i^* \in B_i(a_{-i}^*) \quad \forall i \in N$$

**Lemma 3.1** (pg.19). A strategic game  $\langle N, (A_i), (\succsim_i) \rangle$  has a Nash equilibrium if equivalent to the following statement:

Define set-valued function  $B : A \rightarrow A$  by

$$B(a) = \times_{i \in N} B_i(a_{-i})$$

and there exists  $a^* \in A$  such that  $a^* \in B(a^*)$ .

**Lemma 3.2** (20.1 Kakutani's fixed point theorem). Let  $X$  be a compact convex subset of  $\mathbb{R}^n$  and let  $f : X \rightarrow X$  be a set-valued function for which

- for all  $x \in X$  the set  $f(x)$  is non-empty and convex.
- the graph of  $f$  is closed. (i.e. for all sequences  $\{x_n\}$  and  $\{y_n\}$  such that  $y_n \in f(x_n)$  for all  $n$ ,  $x_n \rightarrow x$  and  $y_n \rightarrow y$  then  $y \in f(x)$ )

Then there exists  $x^* \in X$  such that  $x^* \in f(x^*)$ .

**Definition 3.3** (pg.20). A preference relation  $\succsim_i$  over  $A$  is quasi-concave on  $A_i$  if for every  $a^* \in A$  the upper contour set over  $a_i^*$ , given other players' strategies

$$\{a_i \in A_i : (a_{-i}^*, a_i) \succsim_i a^*\}$$

is convex.

**Proposition 3.1** (20.3). The strategic game  $\langle N, (A_i), (\succsim_i) \rangle$  has a Nash equilibrium if for all  $i \in N$ ,

- the set  $A_i$  of actions of player  $i$  is a nonempty compact convex subset of a Euclidian space

and the preference relation  $\succsim_i$  is

- continuous
- quasi-concave on  $A_i$ .

*Proof.* Let  $B : A \rightarrow A$  be a correspondence defined as

$$B(a) := \times_{i \in N} B_i(a_{-i})$$

Note that for each  $a \in A$  and for each  $i \in N$ ,

$B_i(a_{-i}) \neq \emptyset$  since preference  $\succsim_i$  is continuous and  $A_i$  is compact (EVT).

Also  $B_i(a_{-i})$  is convex since it's basically an intersection of upper contour sets and each of those upper contour is convex since  $\succsim_i$  is quasi-concave.

So the Cartesian product of the finite collection of  $B_i$  is non-empty and convex.

Also the graph  $B$  is closed since  $\succsim_i$  is continuous.

So there exists  $a^* \in A$  such that  $a^* \in B(a^*)$ .

So Nash equilibrium presents. ■

## 4 Lecture 6. Extensive Form Games and Subgame Perfection

### 4.1 Extensive Form Game

**Definition 4.1** (89.1). An **extensive game with perfect information** has the following components.

- A set  $N$  of **players**.
- A set  $H$  of sequences (finite or infinite) of **histories** with properties:
  - $\emptyset \in H$ .
  - For all  $L < K$ ,  $(a^k)_{k=1,2,\dots,K} \in H \implies (a^k)_{k=1,2,\dots,L} \in H$ .
  - For infinite sequence  $(a^k)_{k=1}^\infty$ ,  $(a^k)_{k=1,2,\dots,L} \in H, \forall L \in \mathbb{Z}_{++} \implies (a^k)_{k=1}^\infty \in H$ .

And each component of history  $h \in H$  is an **action** taken by a player.

- A function  $P : H \setminus Z \rightarrow N$ , where for  $h \in H$ ,  $P(h) \in N$  is defined by the player who takes an action after the history  $h$ .
- For each player  $i \in N$  a **preference relation**  $\succsim_i$  defined on  $Z$ .

**Notation 4.1** (pg.90). An extensive game with perfect information can be represented by a 4-tuple,  $\langle N, H, P, (\succsim_i) \rangle$ . *Sometimes it is convenient to specify the structure of an extensive game without specifying the players' preference, as  $\langle N, H, P \rangle$ .*

**Definition 4.2** (pg.90). A history  $(a^k)_{k=1,2,\dots,K} \in H$  is **terminal** if

1. it is infinite,
2. or (i.e. it cannot be extended to another valid history sequence)

$$\forall a^{K+1}, (a^k)_{k=1,2,\dots,K+1} \notin H$$

The set of terminal histories is denoted by  $Z$ .

**Notation 4.2** (pg.90, the action set). After any nonterminal history,  $h \in H \setminus Z$ , the player  $P(h)$  chooses an action from set

$$A(h) = \{a : (h, a) \in H\}$$

**Remark 4.1.** Note that all player function, action set and player preference relation are defined on  $H$ . Thus, unlike a normal form game, which was *player oriented*, we'd better consider an extensive form game as *history oriented*.

**Definition 4.3** (pg.90). We refer to the empty set, which is required to be an element of  $H$ , as the **initial history**.

**Definition 4.4** (92.1). A **strategy of player**  $i \in N$ ,  $s_i$ , in an extensive game with perfect information  $\langle N, H, P, (\succsim_i) \rangle$  is a function that assigns an action in  $A(h)$  to each nonterminal history  $h \in H \setminus Z$  for which  $P(h) = i$ .

**Remark 4.2** (pg.92). A strategy specifies the action chosen by a player for *every* history after which it is his turn to move, *even for histories that is, if the strategy is followed, are never reached*.

**Definition 4.5** (pg.93). For each strategy profile  $s = (s_i)_{i \in N}$  in the extensive game  $\langle N, H, P, (\succsim_i) \rangle$ , the **outcome** of  $s$ ,  $O(s)$ , is defined as the terminal history that results when each player  $i \in N$  follows the precepts of  $s_i$ . That is,  $O(s)$  is the (possibly infinite) history

$$(a^1, \dots, a^K) \in Z$$

such that

$$\forall k \in \{0, 1, \dots, K-1\}, s_{P(a^1, \dots, a^k)}(a^1, \dots, a^k) = a^{k+1}$$

**Definition 4.6** (93.1). A **Nash equilibrium of an extensive game with perfect information**  $\langle N, H, P, (\succsim_i) \rangle$  is a strategy profile  $s^*$  such that for every player  $i \in N$  we have

$$\forall s_i \in S_i, O(s_{-i}^*, s_i^*) \succsim_i O(s_{-i}^*, s_i)$$

**Definition 4.7** (94.1). The **strategic form of the extensive game with perfect information**,  $\Gamma = \langle N, H, P, (\succsim_i) \rangle$ , is the strategic game  $\langle N, (S_i), (\succsim'_i) \rangle$  in which for each player  $i \in N$

- $S_i$  is the **set of strategies** of player  $i$  in  $\Gamma$ .
- $\succsim'_i$  is defined on  $\times_{i \in N} S_i$  and defined by

$$\forall s, s' \in \times_{i \in N} S_i, s \succsim'_i s' \iff O(s) \succsim_i O(s')$$

**Definition 4.8** (pg.94). A **reduced strategy** of player  $i$  is defined to be a function  $f_i$  whose domain is a *subset* of  $\{h \in H : P(h) = i\}$  and has the following properties

1. it associates with every history  $h$  in the domain of  $f_i$  an action in  $A(h)$ .
2. a history  $h$  with  $P(h) = i$  is in the domain of  $f_i$  if and only if all the actions of player  $i$  in  $h$  are those dictated by  $f_i$ . (i.e., for any  $h = (a^k)$  and for any  $h' = (a^k)_{k=1}^L$  as a subsequence of  $h$  such that  $P(h') = i$ ,  $f_i(h') = a^{L+1}$ .)

**Remark 4.3** (pg.94). Each **reduced strategy** of player  $i$  corresponds to a set of strategies of player  $i$ , such that for each vector of strategies of the other players each strategy in this set yields the same outcome. (strategies in the same set are **outcome-equivalent**.)

That's, for each strategy  $s_i \in S_i$ , its reduced strategy can be defined with an outcome equivalence class,  $[s_i]$ ,

$$[s_i] \equiv \{s'_i \in S_i : \forall s_{-i} \in \times_{j \in N \setminus \{i\}} S_j, O(s_{-i}, s_i) = O(s_{-i}, s'_i)\}$$

But in some other game, the definition of outcome-equivalence is more general and defined by generating the same payoff (through possibly difference outcomes), then the reduced strategy is defined as

$$[s_i] \equiv \{s'_i \in S_i : \forall s_{-i} \in \times_{j \in N \setminus \{i\}} S_j, \forall j \in N, O(s_{-i}, s_i) \sim_j O(s_{-i}, s'_i)\}$$

**Definition 4.9** (95.1.1). Let  $\Gamma = \langle N, H, P, (\succsim_i) \rangle$  be an extensive game with perfect information and let  $\langle N, (S_i), (\succsim'_i) \rangle$  be its strategic form. For any  $i \in N$  define the strategies  $s_i, s'_i \in S_i$  to be **equivalent** if

$$\forall s_{-i} \in S_{-i}, \forall j \in N, (s_{-i}, s_i) \sim'_j (s_{-i}, s'_i)$$

**Definition 4.10** (95.1.2). The **reduced strategic form** of  $\Gamma$  is the strategic game  $\langle N, (S'_i), (\succsim''_i) \rangle$  in which for each  $i \in N$  each set  $S'_i$  contains one member of each set of equivalent strategies in  $S_i$  and  $\succsim''_i$  is the preference ordering over  $\times_{j \in N} S'_j$  induced by  $\succsim'_i$ .

## 4.2 Subgame Perfection

**Definition 4.11** (97.1). The **subgame of extensive game with perfect information**  $\Gamma = \langle N, H, P, (\succsim_i) \rangle$  **that follows the history**  $h$  is the extensive game  $\Gamma(h) = \langle N, H|_h, P|_h, (\succsim_i|_h) \rangle$  where

- $H|_h$  is the set of sequences  $h'$  such that  $(h, h') \in H$ .
- $P|_h$  is defined by  $P|_h(h') = P(h, h')$  for each  $h' \in H|_h$ .
- $\succsim_i|_h$  is defined by  $h' \succsim_i|_h h'' \iff (h, h') \succsim_i (h, h'') \in Z$ .

**Notation 4.3** (pg.97). Given strategy  $s_i \in S_i$  and  $h \in H \in \Gamma$ ,  $s_i|_h$  represents the **strategy that  $s_i$  induces in the subgame  $\Gamma(h)$** . That's, for each  $h' \in H|_h$

$$s_i|_h(h') \equiv s_i(h, h')$$

**Notation 4.4.** Let  $O_h$  denote the **outcome function of  $\Gamma(h)$** , that's, for all  $h' \in H|_h$ ,

$$O_h(h') \equiv O(h, h')$$

**Definition 4.12** (97.2). A **subgame perfect equilibrium of an extensive game with perfect information**  $\Gamma = \langle N, H, P, (\succsim_i) \rangle$  is a strategy profile  $s^*$  such that for every player  $i \in N$  and every nonterminal history  $h \in H \setminus Z$  for which  $P(h) = i$  we have

$$O_h(s_{-i}^*|_h, s_i^*|_h) \succsim_i|_h O_h(s_{-i}^*|_h, s_i|_h)$$

for every strategy  $s_i$  of player  $i$  in the subgame  $\Gamma(h)$ .

**Definition 4.13** (pg.97). Equivalently, define SPNE to be a strategy profile  $s^*$  in  $\Gamma$  for which for any history  $h \in H$  the strategy profile  $s^*|_h$  is a Nash equilibrium of the subgame  $\Gamma(h)$ .



**Remark 4.4** (pg. 97). The notion of SPNE requires the action prescribed by each player's strategy to be optimal, given other players' strategies, after *every* history.

**Proposition 4.1** (99.2). Every finite extensive game with perfect information has a subgame perfect equilibrium.