# ECO326 Advanced Microeconomic Theory A Course in Game Theory

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Github Page https://github.com/TianyuDu/Spikey\_UofT\_Notes Note Page TianyuDu.com/notes

Readme this note is based on the course content of ECO326 Advanced Microeconomics - Game Theory, this note contains all materials covered during lectures and mentioned in the course syllabus. However, notations, statements of theorems and proofs are following the book A Course in Game Theory by Osborne and Rubinstein, so they might be, to some extent, more mathematical than the required text for ECO326, An Introduction to Game Theory.

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### 1 Lecture 1. Games and Dominant Strategies

**Assumption 1.1** (pg.4). Assume that each decision-maker is *rational* in the sense that he is aware of his alternatives, forms expectation about any unknowns, has clear preferences, and chooses his action deliberately after some process of optimization.

**Definition 1.1** (pg.4). A model of rational choice consists

- $\bullet$  A set A of actions.
- $\bullet$  A set C of consequences.
- A consequence function  $g: A \to C$ .
- A preference relation  $\succeq$  on C.

**Definition 1.2** (pg.7). A **preference relation** is a complete reflexive and transitive binary relation.

**Definition 1.3** (11.1). A strategic game consists of

- a finite set of **players** N.
- for each player  $i \in N$ , an **actions**  $A_i \neq \emptyset$ .
- for each player  $i \in N$ , a **preference relation**  $\succeq_i$  defined on  $A \equiv \times_{i \in N} A_i$ .

and can be written as a triple  $\langle N, (A_i), (\succeq_i) \rangle$ .

**Definition 1.4** (pg.11). A strategic game  $\langle N, (A_i), (\succsim_i) \rangle$  is **finite** if

$$|A_i| < \aleph_0 \ \forall i \in N$$

## 2 Lecture 2. Iterated Elimination and Rationalizability

#### 2.1 Iterated Elimination of Strictly Dominated Strategies (Actions)

**Definition 2.1** (60.2). The set  $X \subseteq A$  of outcomes of a finite strategic game  $\langle N, (A_i), (u_i) \rangle$  survives iterated elimination of strictly dominated actions if  $X = \times_{j \in N} X_j$  and there is a collection  $\overline{((X_j^t)_{j \in N})_{t=0}^T}$  of sets that satisfies the following conditions for each  $j \in N$ .

- $X_i^0 = A_j$  and  $X_i^T = X_j$ .
- $X_j^{t+1} \subseteq X_j^t$  for each  $t = 0, \dots, T-1$ .
- For each t = 0, ..., T-1 every action of player j in  $X_j^t \setminus X_j^{t+1}$  is strictly dominated in the game  $\langle N, (X_i^t), (u_i^t) \rangle$ , where  $u_i^t$  for each  $i \in N$  is the function  $u_i$  restricted to  $\times_{j \in N} X_i^t$ .
- No action in  $X_t^T$  is strictly dominated in game  $\langle N, (X_i^T), (u_i^T) \rangle$ .

**Proposition 2.1** (61.2). If  $X = \times_{j \in N} X_j$  survives iterated elimination of strictly dominated actions in a <u>finite</u> strategic game  $\langle N, (A_i), (u_i) \rangle$  then  $X_j$  is the set of player j's rationalizable actions for each  $j \in N$ .

#### 2.2 Rationalizability

**Definition 2.2** (pg. 54). A **belief** of player i (about the actions of the other players) is a probability measure,  $\mu_i$ , on  $A_{-i} = \times_{j \in N \setminus \{i\}} A_j$ .

**Definition 2.3** (59.1). An action of player i in a strategic game is a **never** best response if it is not a best response to any belief of player i.

**Definition 2.4** (59.2). The action  $a_i \in A_i$  of player i in the strategic game  $\langle N, (A_i), (u_i) \rangle$  is **strictly dominated** if there is a mixed strategy  $\alpha_i$  of player i such that

$$U_i(a_{-i}, \alpha_i) > u_i(a_{-i}, a_i)$$

for all  $a_{-i} \in A_{-i}$ , where  $U_i(a_{-i}, \alpha_i)$  is the payoff of player i if he uses the mixed strategy  $\alpha_i$  and the other players' vector of actions is  $a_{-i}$ .

## 3 Lecture 3. Nash Equilibrium

**Definition 3.1** (14.1). A Nash equilibrium of a strategic game  $\langle N, (A_i), (\succeq_i) \rangle$  is a profile  $a^* \in A$  of actions with property that for every player  $i \in N$ 

$$(a_i^*, a_{-i}^*) \succsim_i (a_i, a_{-i}^*) \forall a_i \in A_i$$

**Definition 3.2** (pg.15). The **best-response function** for a player i is defined as

$$B_i(a_{-i}) = \{a_i \in A_i : (a_i, a_{-i}) \succeq_i (a'_i, a_{-i}) \ \forall a'_i \in A_i\}$$

**Remark 3.1.** The best-response of  $a_{-i}$  can be written as

$$B_i(a_{-i}) = \bigcap_{a_i' \in A_i} \{ a_i \in A_i : (a_i, a_{-i}) \succsim_i (a_i', a_{-i}) \}$$

where each of them is the upper contour set of  $a'_i$ .

Thus, if  $\succeq_i$  is quasi-concave, then  $B_i(a_{-i})$  is an intersection of convex sets and therefore itself convex.

**Remark 3.2** (pg.15). So a Nash equilibrium is a profile  $a^* \in A$  such that

$$a_i^* \in B_i(a_{-i}^*) \ \forall i \in N$$

**Lemma 3.1** (pg.19). A strategic game  $\langle N, (A_i), (\succsim_i) \rangle$  has a Nash equilibrium if equivalent to the following statement:

Define set-valued function  $B: A \to A$  by

$$B(a) = \times_{i \in N} B_i(a_{-i})$$

and there exists  $a^* \in A$  such that  $a^* \in B(a^*)$ .

**Lemma 3.2** (20.1 Kakutani's fixed point theorem). Let X be a <u>compact</u> convex subset of  $\mathbb{R}^n$  and let  $f: X \to X$  be a set-valued function for which

- for all  $x \in X$  the set f(x) is non-empty and convex.
- the graph of f is closed. (i.e. for all sequences  $\{x_n\}$  and  $\{y_n\}$  such that  $y_n \in f(x_n)$  for all  $n, x_n \to x$  and  $y_n \to y$  then  $y \in f(x)$ )

Then there exists  $x^* \in X$  such that  $x^* \in f(x^*)$ .

**Definition 3.3** (pg.20). A preference relation  $\succeq_i$  over A is quasi-concave on  $A_i$  if for every  $a^* \in A$  the upper contour set over  $a_i^*$ , given other players' strategies

$$\{a_i \in A_i : (a_{-i}^*, a_i) \succeq_i a^*\}$$

is convex.

**Proposition 3.1** (20.3). The strategic game  $\langle N, (A_i), (\succeq_i) \rangle$  has a Nash equilibrium if for all  $i \in N$ ,

• the set  $A_i$  of actions of player i is a nonempty <u>compact convex</u> subset of a Euclidian space

and the preference relation  $\succeq_i$  is

- continuous
- quasi-concave on  $A_i$ .

*Proof.* Let  $B:A\to A$  be a correspondence defined as

$$B(a) := \times_{i \in N} B_i(a_{-i})$$

Note that for each  $a \in A$  and for each  $i \in N$ ,

 $B_i(a_{-i}) \neq \emptyset$  since preference  $\succeq_i$  is continuous and  $A_i$  is compact (EVT).

Also  $B_i(a_{-i})$  is convex since it's basically an intersection of upper contour sets and each of those upper contour is convex since  $\succeq_i$  is quasi-concave. So the Cartesian product of the finite collection of  $B_i$  is non-empty and

So the Cartesian product of the finite collection of  $B_i$  is non-empty and convex.

Also the graph B is closed since  $\succsim_i$  is continuous.

So there exists  $a^* \in A$  such that  $a^* \in B(a^*)$ .

So Nash equilibrium presents.

## 4 Lecture 6. Extensive Form Games and Subgame Perfection

#### 4.1 Extensive Form Game

**Definition 4.1** (89.1). An extensive game with perfect information has the following components.

- $\bullet$  A set N of players.
- A set *H* of sequences (finite or infinite) of **histories** with properties:
  - $-\emptyset \in H$ .
  - For all L < K,  $(a^k)_{k=1,2,...,K} \in H \implies (a^k)_{k=1,2,...,L} \in H$ .
  - For infinite sequence  $(a^k)_{k=1}^{\infty}$ ,  $(a^k)_{k=1,2,\dots,L} \in H, \ \forall L \in \mathbb{Z}_{++} \implies (a^k)_{k=1}^{\infty} \in H.$

And each component of history  $h \in H$  is an **action** taken by a player.

- A function  $P: H \setminus Z \to N$ , where for  $h \in H$ ,  $P(h) \in N$  is defined by the player who takes an action after the history h.
- For each player  $i \in N$  a **preference relation**  $\succsim_i \underline{\text{defined on } Z}$ .

**Notation 4.1** (pg.90). An extensive game with perfect information can be represented by a 4-tuple,  $\langle N, H, P, (\succeq_i) \rangle$ . Sometimes it is convenient to specify the structure of an extensive game without specifying the players' preference, as  $\langle N, H, P \rangle$ .

**Definition 4.2** (pg.90). A history  $(a^k)_{k=1,2,...,K} \in H$  is terminal if

- 1. it is infinite,
- 2. or (i.e. it cannot be extended to another valid history sequence)

$$\forall a^{K+1}, \ (a^k)_{k=1,2,...,K+1} \notin H$$

The set of terminal histories is denoted by Z.

**Notation 4.2** (pg.90, the action set). After any nonterminal history,  $h \in H \setminus Z$ , the player P(h) chooses an action from set

$$A(h) = \{a : (h, a) \in H\}$$

**Remark 4.1.** Note that all player function, action set and player preference relation are defined on H. Thus, unlike a normal form game, which was player oriented, we'd better consider an extensive form game as history oriented.

**Definition 4.3** (pg.90). We refer to the empty set, which is required to be an element of H, as the **initial history**.

**Definition 4.4** (92.1). A strategy of player  $i \in N$ ,  $s_i$ , in an extensive game with perfect information  $\langle N, H, P, (\succeq_i) \rangle$  is a function that assigns an action in A(h) to each nonterminal history  $h \in H \setminus Z$  for which P(h) = i.

Remark 4.2 (pg.92). A strategy specifies the action chosen by a player for every history after which it is his turn to move, even for histories that is, if the strategy is followed, are never reached.

**Definition 4.5** (pg.93). For each strategy profile  $s = (s_i)_{i \in N}$  in the extensive game  $\langle N, H, P, (\succeq_i) \rangle$ , the **outcome** of s, O(s), is defined as the <u>terminal history</u> that results when each player  $i \in N$  follows the precepts of  $s_i$ . That is, O(s) is the (possibly infinite) history

$$(a^1,\ldots,a^K)\in Z$$

such that

$$\forall k \in \{0, 1, \dots K - 1\}, \ s_{P(a^1, \dots, a^k)}(a^1, \dots, a^k) = a^{k+1}$$

#### 4.2 Subgame Perfection