

STA447: Stochastic Processes

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January 19, 2020

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1 Markov Chain Probabilities

Definition 1.1. A **discrete-time, discrete-space, and time-homogenous Markov chain** is a triple of (S, v, p) in which

- (i) S represents the *state space*, which is nonempty and countable;
- (ii) *initial probability* v , which is a distribution on S ;
- (iii) and *transition probability* p_{ij} .

Definition 1.2. A Markov chain satisfies the **time-homogenous property** if

$$P(X_{n+1} = j | X_n = i) = P(X_1 = j | X_0 = i) = p_{ij} \quad \forall n \in \mathbb{N} \quad (1.1)$$

Definition 1.3. A Markov chain satisfies the **Markov property** if

$$P(X_{n+1} = j | X_n = i_n, X_{n-1} = i_{n-1}, \dots, X_0 = i_0) = P(X_{n+1} = j | X_n = i_n) \quad (1.2)$$

That is, the chain is *memoryless*.

Proposition 1.1 (Multistep Arrival Probability). Let $m = |S|$ and $\mu_i^{(n)} := P(X_n = i)$ denote the probability that the state ends up at i after n step. By the law of total expectation,

$$P(X_n = i) = \sum_{j \in S} P(X_n = i, X_{n-1} = j) \quad (1.3)$$

$$= \sum_{j \in S} P(X_n = i | X_{n-1} = j) P(X_{n-1} = j) \quad (1.4)$$

$$= \sum_{j \in S} P(X_{n-1} = j) p_{ij} \quad (1.5)$$

$$= \sum_{j \in S} \mu_j^{(n-1)} p_{ij} \quad (1.6)$$

Let $\mu^{(n)} := [\mu_1^{(n)}, \mu_2^{(n)}, \dots, \mu_m^{(n)}] \in \mathbb{R}^{1 \times m}$ and $P = [p_{ij}] \in \mathbb{R}^{m \times m}$. In matrix notation:

$$\mu^{(n)} = \mu^{(n-1)} P \quad (1.7)$$

where $\mu^{(0)} = v = [v_1, v_2, \dots, v_m]$. Define $P^0 = I_m$, then

$$\mu^{(n)} = v P^n \quad (1.8)$$

Proposition 1.2 (Multistep Transition Probability). Define $p_{ij}^{(n)} := P(X_{m+n} = j | X_m = i)$ to be the probability of arriving state j after n steps, starting from state i . By the time-homogenous property,

$$p_{ij}^{(n)} = P(X_{m+n} = j | X_m = i) \quad \forall m \in \mathbb{N} \quad (1.9)$$

Let $P^{(n)} := [p_{ij}^{(n)}] \in \mathbb{R}^{m \times m}$.

Initial Step: for $n = 1$, $P^{(1)} = P$ by definition.

Inductive Step: for $n \in \mathbb{N}$,

$$p_{ij}^{(n+1)} = P(X_{n+1} = j | X_0 = i) \quad (1.10)$$

$$= \sum_{k \in S} P(X_{n+1} = j | X_n = k, X_0 = i) P(X_n = k | X_0 = i) \quad (1.11)$$

$$= \sum_{k \in S} P(X_{n+1} = j | X_n = k) p_{ik}^{(n)} \quad (1.12)$$

$$= \sum_{k \in S} p_{ik}^{(n)} p_{kj} \quad (1.13)$$

$$= [P^{(n)} P]_{ij} \quad (1.14)$$

Therefore,

$$P^{(n+1)} = P^{(n)} P \quad (1.15)$$

and

$$P^{(n)} = P^n \quad (1.16)$$

Theorem 1.1 (Chapman-Kolmogorov Equations). Let $n = (n_1, n_2, \dots, n_k)$ be a multi-set of non-negative integers, then

$$P^{(\sum_{i=1}^k n_i)} = \prod_{i=1}^k P^{(n_i)} \quad (\dagger) \quad (1.17)$$

Proof. Prove by induction on the size of multi-set:

Base case is trivial for $k = 1$.

Inductive step for $k > 1$, suppose (\dagger) holds for every set of length k , consider another multi-set with length $k + 1$: $n' = (n_1, n_2, \dots, n_k, n_{k+1})$. Let $\delta := \sum_{i=1}^k n_i$.

$$P_{ij}^{(\delta+n_{k+1})} = P(X_{\delta+n_{k+1}} = j | X_0 = i) \quad (1.18)$$

$$= \sum_{k \in S} P(X_{\delta+n_{k+1}} = j | X_\delta = k, X_0 = i) P(X_\delta = k | X_0 = i) \quad (1.19)$$

$$= \sum_{k \in S} P(X_{\delta+n_{k+1}} = j | X_\delta = k) P(X_\delta = k | X_0 = i) \quad (1.20)$$

$$= \sum_{k \in S} P(X_{n_{k+1}} = j | X_0 = k) P(X_\delta = k | X_0 = i) \quad (1.21)$$

$$= \sum_{k \in S} p_{kj}^{n_{k+1}} p_{ik}^{(\delta)} \quad (1.22)$$

$$= [P^{(\delta)} P^{(n_{k+1})}]_{ij} \quad (1.23)$$

$$\implies P^{(\delta+n_{k+1})} = P^{(\delta)} P^{(n_{k+1})} \quad (1.24)$$

■

Corollary 1.1 (Chapman-Kolmogorov Inequality). For every $k \in S$,

$$p_{ij}^{(m+n)} \geq p_{ik}^{(m)} p_{kj}^{(n)} \quad (1.25)$$

For $k, \ell \in S$,

$$p_{ij}^{(m+s+n)} \geq p_{ik}^{(m)} p_{k\ell}^{(s)} p_{\ell j}^{(n)} \quad (1.26)$$

Notation 1.1. Let $N(i) := |\{n \geq 1 : X_n = i\}|$ denote the number of arrivals to state i of the chain.

Definition 1.4. Define the **return probability** from state i to j , f_{ij} , as the probability of arriving state j starting from state i . That is,

$$f_{ij} = P(\exists n \geq 1 \text{ s.t. } X_n = j | X_0 = i) \quad (1.27)$$

$$= P(N(j) \geq 1 | X_0 = i) \quad (1.28)$$

Proposition 1.3. The probability of firstly arriving j , then arriving k (denoted as event E) starting from i equals

$$P(E | X_0 = i) = f_{ij} f_{jk} \quad (1.29)$$

Proof. The proof follows the time-homogenous property. ■

Corollary 1.2.

$$P(N(i) \geq k | X_0 = i) = (f_{ii})^k \quad (1.30)$$

$$P(N(j) \geq k | X_0 = i) = f_{ij} (f_{jj})^{k-1} \quad (1.31)$$

Definition 1.5. A state i in a Markov chain is **recurrent** if $f_{ii} = 1$. Otherwise, this state is **transient**.

Theorem 1.2 (Recurrent State Theorem). The following statements are equivalent:

- (i) State i is recurrent (i.e., $f_{ii} = 1$);
- (ii) $P(N(i) = \infty | X_0 = i) = 1$, that is, state i will be visited infinitely often;
- (iii) $\sum_{n=1}^{\infty} p_{ii}^{(n)} = \infty$.

The following statements are equivalent:

- (a) State i is transient;
- (b) $P(N(i) = \infty | X_0 = i) = 0$, that is, state i will only be visited finitely many times;
- (c) $\sum_{n=1}^{\infty} p_{ii}^{(n)} < \infty$.

Proof. We only show the equivalence of (i) \sim (iii), (a) \sim (c) are simply the negation of previous statements.

(i) \iff (ii):

$$P(N(i) = \infty | X_0 = i) = P(\lim_{k \rightarrow \infty} N(i) \geq k | X_0 = i) \quad (1.32)$$

$$= \lim_{k \rightarrow \infty} P(N(i) \geq k | X_0 = i) \quad (1.33)$$

$$= \lim_{k \rightarrow \infty} (f_{ii})^k = 1 \text{ if and only if } f_{ii} = 1 \quad (1.34)$$

(i) \iff (iii):

$$\sum_{n=1}^{\infty} p_{ii}^{(n)} = \sum_{n=1}^{\infty} P(X_n = i | X_0 = i) \quad (1.35)$$

$$= \sum_{n=1}^{\infty} \mathbb{E}(1_{X_n=i} | X_0 = i) \quad (1.36)$$

$$= \mathbb{E} \left(\sum_{n=1}^{\infty} 1_{X_n=i} \middle| X_0 = i \right) \quad (1.37)$$

$$= \mathbb{E}(N(i) | X_0 = i) \quad (1.38)$$

$$= \sum_{n=k}^{\infty} k P(N(i) = k | X_0 = i) \quad (1.39)$$

$$= \sum_{n=k}^{\infty} P(N(i) \geq k | X_0 = i) \quad (1.40)$$

$$= \sum_{n=k}^{\infty} (f_{ii})^k \quad (1.41)$$

$$= \infty \text{ if and only if } f_{ii} = 1 \quad (1.42)$$

■

Lemma 1.1 (Stirling's Approximation).

$$n! \approx (n/e)^n \sqrt{2\pi n} \quad (1.43)$$

Proposition 1.4. For simple random walk, if $p = 1/2$, then $f_{ii} = 1 \ \forall i \in S$. Otherwise, all states are transient.

$$\forall i \in S, \ f_{ii} = 1 \iff p = \frac{1}{2} \quad (1.44)$$

Proof. For simplicity, consider state 0 and the series $\sum_{n=1}^{\infty} p_{00}^{(n)}$. Note that for odd n 's, $p_{00}^{(n)} = 0$.

For all even n 's such that $n = 2k$,

$$\sum_{n=1}^{\infty} p_{00}^{(n)} = \sum_{k=1}^{\infty} p_{00}^{(2k)} \quad (1.45)$$

$$= \sum_{k=1}^{\infty} \binom{2k}{k} p^k (1-p)^k \quad (1.46)$$

$$= \sum_{k=1}^{\infty} \frac{2k!}{(k!)^2} p^k (1-p)^k \quad (1.47)$$

$$\approx \sum_{k=1}^{\infty} \frac{(2k/e)^{2k} \sqrt{4\pi k}}{(k^k e^{-k} \sqrt{2\pi k})^2} p^k (1-p)^k \quad (1.48)$$

$$= \sum_{k=1}^{\infty} \frac{2^{2k} k^{2k} e^{-2k} 2\sqrt{\pi k}}{k^{2k} e^{-2k} 2\pi k} p^k (1-p)^k \quad (1.49)$$

$$= \sum_{k=1}^{\infty} \frac{2^{2k}}{\sqrt{\pi k}} p^k (1-p)^k \quad (1.50)$$

$$= \sum_{k=1}^{\infty} \frac{4^k}{\sqrt{\pi k}} p^k (1-p)^k \quad (1.51)$$

$$= \sum_{k=1}^{\infty} \frac{1}{\sqrt{\pi k}} [4p(1-p)]^k \quad (1.52)$$

When $p = \frac{1}{2}$,

$$\sum_{n=1}^{\infty} p_{00}^{(n)} = \frac{1}{\sqrt{\pi}} \sum_{k=1}^{\infty} k^{-1/2} \quad (1.53)$$

$$= \infty \quad (1.54)$$

When $p \neq \frac{1}{2}$,

$$\sum_{n=1}^{\infty} p_{00}^{(n)} < \frac{1}{\sqrt{\pi}} \sum_{k=1}^{\infty} [4\pi(1-p)]^k \quad (1.55)$$

$$< \infty \quad (1.56)$$

By the recurrent state theorem, $f_{ii} = 1 \iff p = 1/2$.

For other $i \neq 0$, the prove is similar. ■

Theorem 1.3 (f-Expansion).

$$f_{ij} = p_{ij} + \sum_{k \in S \setminus \{j\}} p_{ik} f_{kj} \quad (1.57)$$

Proof.

$$f_{ij} = P(\exists n \in \mathbb{Z}_{++} \text{ s.t. } X_n = j | X_0 = i) \quad (1.58)$$

$$= \sum_{k \in S} P(\exists n \in \mathbb{Z}_{++} \text{ s.t. } X_n = j | X_0 = i, X_1 = k) P(X_1 = k | X_0 = i) \quad (1.59)$$

$$= \sum_{k \in S} P(\exists n \in \mathbb{Z}_{++} \text{ s.t. } X_n = j | X_1 = k) P(X_1 = k | X_0 = i) \text{ (Markov Property)} \quad (1.60)$$

$$= \underbrace{P(\exists n \in \mathbb{Z}_{++} \text{ s.t. } X_n = j | X_1 = j)}_{=1} P(X_1 = j | X_0 = i) + \sum_{k \neq j} f_{kj} P(X_1 = k | X_0 = i) \quad (1.61)$$

$$= p_{ij} + \sum_{k \neq j} f_{kj} p_{ik} \quad (1.62)$$

■

Definition 1.6. State i is said to **communicate** with state j , denoted as $i \rightarrow j$, if $f_{ij} > 0$.

Proposition 1.5 (Alternative Definition). The following statements are equivalent:

- (i) $i \rightarrow j$;
- (ii) $\exists m \geq 1, \text{ s.t. } p_{ij}^{(m)} > 0$.

Proof. **TODO:** *Proof.*

■

Definition 1.7. A Markov chain is **irreducible** if $i \rightarrow j \forall i, j \in S$.

1.1 Recurrence and Transience Equivalence Theorem

Theorem 1.4 (Sum Lemma). If

- (i) $i \rightarrow k$;
 - (ii) $\ell \rightarrow j$;
 - (iii) $\sum_{n=1}^{\infty} p_{k\ell}^{(n)} = \infty$.
- Then, $\sum_{n=1}^{\infty} p_{ij}^{(n)} = \infty$.

Proof. Suppose $i \rightarrow k$ and $\ell \rightarrow j$, then there exists m and r such that $p_{ik}^{(m)} > 0$ and $p_{\ell j}^{(r)} > 0$.

By the Chapman-Kolmogorov inequality, $p_{ij}^{(m+n+r)} \geq p_{ik}^{(m)} p_{k\ell}^{(n)} p_{\ell j}^{(r)}$.

Then,

$$\sum_{n=1}^{\infty} p_{ij}^{(n)} \geq \sum_{n=m+r+1}^{\infty} p_{ij}^{(n)} \quad (1.63)$$

$$= \sum_{s=1}^{\infty} p_{ij}^{(m+s+r)} \quad (1.64)$$

$$\geq \sum_{s=1}^{\infty} p_{ik}^{(m)} p_{k\ell}^{(s)} p_{\ell j}^{(r)} \quad (1.65)$$

$$= p_{ik}^{(m)} p_{\ell j}^{(r)} \sum_{s=1}^{\infty} p_{k\ell}^{(s)} = \infty \quad (1.66)$$

■

Theorem 1.5. If $i \leftrightarrow k$, then

$$f_{ii} = 1 \iff f_{kk} = 1 \quad (1.67)$$

Proof. **TODO:** *Proof.*

■

Theorem 1.6 (Case Theorem). For an *irreducible* Markov chain, it is either

- (a) a **recurrent** Markov chain: $\forall i \in S, f_{ii} = 1$ and $\sum_{n=1}^{\infty} p_{ij}^{(n)} = \infty \forall i, j \in S$;
- (b) or a **transient** Markov chain: $\forall i \in S, f_{ii} < 1$ and $\sum_{n=1}^{\infty} p_{ij}^{(n)} < \infty \forall i, j \in S$.

Proof. **TODO:** *Proof.*

■

Theorem 1.7 (Finite Space Theorem). An *irreducible* Markov chain on a *finite* state space is always recurrent.

Proof. Let $i \in S$ (u.i.),

$$\sum_{j \in S} \sum_{n=1}^{\infty} p_{ij}^{(n)} = \sum_{n=1}^{\infty} \sum_{j \in S} p_{ij}^{(n)} \quad (1.68)$$

$$= \sum_{n=1}^{\infty} 1 = \infty \quad (1.69)$$

Because S is finite, $\exists k \in S$ such that $\sum_{n=1}^{\infty} p_{ik}^{(n)} = \infty$. Therefore, all states are recurrent.

■

Theorem 1.8 (Hit-Lemma). Define H_{ij} as the event in which the chain starts from j and visits i without firstly returning to j (*direct path from j to i*)¹:

$$H_{ij} := \{\exists n \in \mathbb{N} \text{ s.t. } X_n = i \wedge X_m \neq j \forall m < n\} \quad (1.70)$$

If $j \rightarrow i$ with $j \neq i$, then $P(H_{ij}|X_0 = j) > 0$.

Theorem 1.9 (f-Lemma). For all $i, j \in S$, if $j \rightarrow i$ and $f_{jj} = 1$, then $f_{ij} = 1$.

Proof. For $i = j$, trivial.

Suppose $i \neq j$, since $j \rightarrow i$, then $P(H_{ij}|X_0 = j) > 0$.

Further,

$$P(X_n \neq j \forall n \in \mathbb{Z}_{++} | X_0 = j) \geq P(H_{ij}|X_0 = j)P(X_n \neq j \forall n \in \mathbb{Z}_{++} | X_0 = i) \quad (1.71)$$

$$\implies 0 = 1 - f_{jj} \geq P(H_{ij}|X_0 = j)(1 - f_{ij}) \quad (1.72)$$

$$\implies f_{ij} = 1 \quad (1.73)$$

■

¹Notation abuse: H_{ij} describes the event starting from j and ending at i , instead of the other way round.

Theorem 1.10 (Infinite Returns Lemma). For an *irreducible* Markov chain,

- (i) if this chain is recurrent, then $P(N(j) = \infty | X_0 = i) = 1 \ \forall i, j \in S$;
- (ii) if this chain is transient, then $P(N(j) = \infty | X_0 = i) = 0 \ \forall i, j \in S$.

Proof. Let $i, j \in S$.

Suppose the chain is irreducible and recurrent, if $i = j$, then $f_{ii} = f_{jj} = 1$.

Otherwise, $i \neq j$. Since $j \rightarrow i$, by the f-Lemma, $f_{jj} = f_{ii} = f_{ij} = f_{ji} = 1$.

$$P(N(j) = \infty | X_0 = i) = \lim_{k \rightarrow \infty} P(N(j) \geq k | X_0 = i) \quad (1.74)$$

$$= \lim_{k \rightarrow \infty} f_{ij} f_{jj}^{k-1} \quad (1.75)$$

$$= 1 \quad (1.76)$$

When the chain is transient, $f_{jj} < 1$, and $\lim_{k \rightarrow \infty} f_{ij} f_{jj}^{k-1} = 0$. ■

Theorem 1.11 (Recurrent Equivalences Theorem). For a irreducible Markov chain (so that $i \rightarrow j$ for all $i, j \in S$), the following statements are equivalent:

- (1) $\exists k, \ell \in S$ such that $\sum_{n=1}^{\infty} p_{k\ell}^{(n)} = \infty$;
- (2) $\forall i, j \in S$, $\sum_{n=1}^{\infty} p_{ij}^{(n)} = \infty$;
- (3) $\exists k \in S$ s.t. $f_{kk} = 1$;
- (4) $\forall j \in S$, $f_{jj} = 1$;
- (5) $\forall i, j \in S$, $f_{ij} = 1$;
- (6) $\exists k, \ell \in S$ such that $P(N(\ell) = \infty | X_0 = k) = 1$;
- (7) $\forall i, j \in S$, $P(N(j) = \infty | X_0 = i)$.