## Discrete Mathematics Recitation Class

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## Contents

#### The Natural Numbers

## Propositional Logic

Compound Expressions

Equivalence

Arguments

## Predicate Logic

**Predicates** 

Logical Quantifiers

## Naive Set Theory

Sets

Operations on Sets

**Ordered Pairs** 

Problems in Naive Set Theory

# The Natural Numbers (P8)

The set of natural numbers

$$\mathbb{N}:=\{0,1,2,3,\cdots\}$$

Addition & Multiplication

$$+: \mathbb{N} \times \mathbb{N} \to \mathbb{N} \qquad \cdot: \mathbb{N} \times \mathbb{N} \to \mathbb{N}$$

If  $a, b, c \in \mathbb{N}$ ,

Comutativity : 
$$a + b = b + a$$
,  $a \cdot b = b \cdot a$ 

Associativity : 
$$a + b + c = a + (b + c)$$
,  $a \cdot b \cdot c = a \cdot (b \cdot c)$ 

Distributivity : 
$$a \cdot (b + c) = a \cdot c + b \cdot c$$

# The Natural Numbers (P9)

#### **Definitions**

- 1. greater than (or equal to)
- 2. exact division
- 3. odd & even
- 4. prime

#### Question

Prove that every natural number is either even or odd and not both.

# Uniqueness of Odevity (P9)

#### Proof.

Suppose that  $n \in \mathbb{N}$ ,

- 1. prove that if n is odd, n cannot be even
- 2. prove that if n is even, n cannot be odd
- 3. prove that n cannot be neither odd nor even (Do not forget!)

#### Proof.

Suppose that  $n \in \mathbb{N}$ ,

- 1. prove that if n isnot odd, n must be even
- prove that if n isnot even, n must be odd
- prove that n cannot be both odd and even (Do not forget!)

## **Propositions**

## **Definition** (P11)

- Declarative
- ▶ One can tell true (T) or false (F) immediately

Compound Expressions (using connectives) (P12)

- ▶ unary connective: ¬(not)
- binary connectives: ∨(disjunction), ∧(conjunction),
   ⇒(implication), ⇔(biconditonal)

# Negation (P13)

**Definition** (by truth table) Suppose *A* is a proposition:

$$\begin{array}{c|c} A & \neg A \\ \hline T & F \\ \hline F & T \\ \end{array}$$

# Conjunction (P13)

**Definition** (by truth table) Suppose A, B are propositions:

Α	В	A∧B	
Т	Т	Т	
Т	F	F	
F	Т	F	
F	F	F	

# Disjunction (P14)

**Definition** (by truth table) Suppose A, B are propositions:

Α	В	A∨B	
Т	Т	Т	
Т	F	Т	
F	Т	Т	
F	F	F	

# Implication (P16)

**Definition** (by truth table) Suppose A, B are propositions:

Α	В	A⇒B	
Т	Т	Т	
Т	F	F	
F	Т	Т	
F	F	Т	

A: antecedent, B: consequence

# Biconditional (P18)

**Definition** (by truth table) Suppose A, B are propositions:

Α	В	A⇔B	
Т	Т	Т	
Т	F	F	
F	Т	F	
F	F	Т	

# Logical Equivalence

#### **Definitions**

- 1. tautology (P15)
  - A compound expression that is always true.
- 2. contradiction (P15)
  - A compound expression that is always false.
- 3. logical equivalence (P19)
  - The biconditional of two propositions is tautology.

# Logical Equivalence Examples

1. De Morgan Rules (P19)

$$\neg (A \lor B) \equiv (\neg A) \land (\neg B) \qquad \neg (A \land B) \equiv (\neg A) \lor (\neg B)$$

2. Contraposition (P20)

Α	В	¬A	¬B	A⇒B	¬B⇒¬A	$A \Rightarrow B \Leftrightarrow \neg B \Rightarrow \neg A$
Т	Т	F	F	Т	Т	Т
Т	F	F	Т	F	F	Т
F	Т	Т	F	Т	Т	Т
F	F	Т	Т	T	Т	Т

3. Other Logical Equivalences (P21-P23)

# Manipulating propositional expressions

## **Usages**

- 1. Express some connectives with other connectives (P24) **e.g.** Prove that  $\neg(A \Rightarrow B)$  is equivalent to  $A \land \neg B$
- 2. Prove tautologies/contradictions without truth tables (P25) **e.g.** Prove that  $(A \land B) \Rightarrow (A \lor B)$  is a tautology.

## **Arguments**

## **Definitions** (P27)

- 1. arguments
- 2. premises
- 3. conclusion
- 4. valid

## Rules of Inference (P29-P33)

- 1. Hypothetical Syllogisms
- 2. Disjunctive and Conjunctive Syllogisms
- 3. Simple Arguments

# Validity & Soundness (P34)

#### **Definitions**

- 1. *validity*: True premises leads to true conclusion. (Conclusion does not need to be true if premises are not true.)
- 2. *soundness*: All premises are true. (Only deals with truth of premises.)

## Predicate

## **Definitions** (P38)

- 1. predicate: Declarative sentence involving variables.
- 2. arity: number of distinct variables in the predicate.

Difference between Predicates and Propositions:

#### e.g.

- x > 5 is not a proposition. (we cannot tell the true or false immediately because the value of x is unknown)
- "x < 5" is a unary predicate.



## **Predicate**

```
(P39)
```

 $\frac{\text{bound with quantifiers}}{\text{replace variables by constants}} \text{Sentences}$ 

## Logical Quantifiers

## **Definitions** (P41)

- 1. ∀: for all···
- 2. ∃: there exists···

## Hanging Quantifiers (P42)

Appearing at the back is equivalent to appearing just before the expression.

Contraposition of Quantifiers (P43)

$$\neg((\exists x \in M)A(x)) \equiv (\forall x \in M)\neg A(x)$$

$$\neg((\forall x \in M)A(x)) \equiv (\exists x \in M)\neg A(x)$$

Passing negation symbol in quantifier statements from the very left to the very right

# Vacuous Truth (P44)

#### **Definition**

- ▶ Apply only for universal quantifier ∀
- ▶ Domain  $M = \emptyset$
- ▶ Regardless of A(x) being true or false

A universal statement is true unless there is a counterexample to prove it false.

# Quantifier Order (P45)

### e.g.

- ∀x∃y(x + y > 0) is a true statement.
  For each x, one needs to find a specific y to satisfy the statement, y can be different for different x.
- ∃x∀y(x + y > 0) is a false statement.
   Find a y first, then that specific y (value unchanged) needs to satisfy the statement for any value of x.

# Tautology in Predicate Logic

## **Definition** (P46)

non-empty domain of discourse

## Rules of Inference (P47)

- 1. Universal Instantiation
  - 2. Universal Generalization
    - 3. Existential Instantiation
    - 4. Existential Generalization
- ▶ All other rules of inference that are valid in propositional logic

# Proof by Contradiction

Recall the definition of validity of arguments: If all premises are true, then the conclusion is also true.

#### Question

Prove that there are infinitely many prime numbers.

### Proof.

Suppose that there are not infinitely many primes. Therefore there must be finitely many primes:  $p_1 < p_2 < \cdots < p_n$  and hence a largest prime  $p_n$ . Let

$$m = \prod_{1 \leqslant i \leqslant n} p_i + 1$$

By our assumption, m is not prime. So there exists  $p_j$  such that  $p_j|m$ , i.e.  $m=p_jk$  for some integer k. But now,

$$1 = p_j \left( k - \prod_{1 \leqslant i \leqslant n} p_i \right)$$

which shows that  $p_j|1$ . This is a contradiction. Therefore, we can conclude that there are infinitely many primes.

## Sets

## **Definition** (P53)

- collection of objects
- ignore order
- ignore repetitions

Express Sets with Predicates:  $X = \{x | P(x)\}$  *i.e.*  $x \in X$  iff P(x) **e.g.** 

Suppose  $A = \{x | A(x)\}, B = \{x | B(x)\}, C$  are sets,

- $\forall x \in A(A(x) \land B(x)) \Leftrightarrow A \subseteq B \Leftrightarrow (A(x) \Rightarrow B(x))$
- $\exists x \in A(A(x) \land B(x)) \Leftrightarrow A \cap B \neq \emptyset$
- $C = \{x | \neg A(x)\} = A^c$

The examples above may help you understand the relation between sets and predicates better.

Sets

## Definitions

Operations on Sets

- 1. equality of two sets (P54)
- 2. empty set (P55)
- 3. (proper) subset (P56)
- 4. cardinality (P58)
- 5. powerset (P58)
- Empty set (∅) is subset of any set.
- $A \subseteq X \equiv A \in \mathcal{P}(X)$
- ▶  $|\mathcal{P}(X)| = 2^{|X|}$

**e.g.** (P57)

Let  $A = \{\emptyset, \{\{\emptyset\}\}\}\}$ ,  $B = \{\emptyset\}$  and  $C = \{\{\emptyset\}\}\}$ .  $B \subseteq A$ , but  $B \notin A$ 

A, and  $C \in A$ , but  $C \nsubseteq A$ .



# Operations on Sets (P59-P60)

## Definitions

- 1. union
- 2. intersection
- 3. difference
- 4. complement
- 5. disjoint

# Complex Operations on Sets Suppose A, B, C are sets,

- $ightharpoonup A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$
- $A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$
- $(A \cup B) \setminus C = (A \setminus C) \cup (B \setminus C)$
- $(A \cap B) \setminus C = (A \setminus C) \cap (B \setminus C)$
- $ightharpoonup A \setminus (B \cup C) = (A \setminus C) \cap (B \setminus C)$
- $A \setminus (B \cap C) = (A \setminus C) \cup (B \setminus C)$
- $\triangleright$   $A \setminus B = B^c \cap A$
- $(A \setminus B)^c = A^c \cup B$

Use Wenn Diagrams to help understand.

Sets

# Operations on Sets (P61-P62)

**e.g.** For a finite number  $n \in \mathbb{N}$  of sets  $A_0, A_1, \dots, A_n$ :

$$\bigcup_{k=0}^{n} A_k := A_0 \cup A_1 \cup \cdots \cup A_n \qquad \bigcap_{k=0}^{n} A_k := A_0 \cap A_1 \cap \cdots \cap A_n$$

If *n* extends to  $\infty$ ,

$$x \in \bigcup_{k=0}^{\infty} A_k : \Leftrightarrow \underset{k \in \mathbb{N}}{\exists} x \in A_k \qquad x \in \bigcap_{k=0}^{\infty} A_k : \Leftrightarrow \underset{k \in \mathbb{N}}{\exists} x \in A_k$$

# Operations on Sets (P63)

If A is an set and  $X \subseteq \mathcal{P}(A)$ , then

$$\bigcup X = \{x \in A | (\exists y \in X)(x \in y)\}\$$

and

Sets

$$\bigcap X = \{x \in A | (\forall y \in X)(x \in y)\}$$

**e.g.** Let  $A = \{1, 2\}, \mathcal{P}(A) = \{\emptyset, \{1\}, \{2\}, \{1, 2\}\}$ . Suppose  $X = \{\{1\}, \{1, 2\}\}$ , then

$$\bigcap X = \{1\}, \bigcup X = \{1, 2\}$$

Here  $\bigcup X$  indicates the union of all objects (sets) in X, while  $\bigcap X$  indicates the intersection of all objects (set) in X.

e.g. Let

$$X = \{A \in \mathcal{P}(\mathbb{N}) | (\exists k \in \mathbb{N}) (\forall n \in \mathbb{N}) (n \in A \lor n = k) \}$$

Then

$$\bigcup X = \mathbb{N} \text{ and } \bigcap X = \emptyset$$

Sets

# Ordered Pairs & Cartesian Products (P64-P65)

#### **Definitions**

- 1.  $(a,b) := \{\{a\}, \{a,b\}\}$
- $2. (a,b) = (c,d) \Leftrightarrow (a=c) \land (b=d)$
- 3. Cartesian product of two sets  $A \times B := \{(a, b) | a \in A \land b \in B\}.$
- 4. ordered n—tuple & n—fold Cartesian product

# Paradoxes (P67-P69)

- Epimenides Paradox
- ► Barber Paradox (Russell's Paradox)

#### Russell's Paradox:

#### **Theorem**

The set of all sets that are not members of themselves is not a set. i.e.

$$R := \{x | x \notin x\}$$
 is not a set.

#### Proof.

The proof is by contradiction. Suppose that R is a set. If  $R \in R$ , then  $R \notin R$  by the definition of R, which is a contradiction. If  $R \notin R$ , then  $R \in R$  by the definition of R, which is also a contradiction.