

Ve 280

Programming and Introductory Data Structures

ADT Example: IntSet;
Improve ADT Efficiency;

Outline

- A Class Example: a Mutable Set of Integers (IntSet)
- Improve the Efficiency of IntSet

Abstract Data Types

Using Classes

- Next, consider the three final routines:
 - query: search the array looking for a specific number.
 - remove: search the array for a number; if it exists, remove it.
 - insert: search the array for a number; if it doesn't exist, add it.
- All three of these have "search" in common.
- One might be tempted to just write `insert` and `remove` in terms of `query`, will this work?
 - Hint: think about `remove`.
- `query` only tells us **whether** the element exists, not **where** – we need one more method...

Abstract Data Types

Using Classes

```
const int MAXELTS = 100;
class IntSet { // OVERVIEW omitted for space
    int      elts[MAXELTS];
    int      numElts;
```

```
    int indexOf(int v);
        // EFFECTS: returns the index of
        //           v if it exists in the
        //           array, MAXELTS otherwise.
```

```
public:
    void insert(int v);
    void remove(int v);
    bool query(int v);
    int  size();
};
```

Note: This member function must be **private**. This is because it exposes details about the concrete representation. It is inappropriate to expose these details to a user of this class.

Abstract Data Types

Using Classes

```
const int MAXELTS = 100;
class IntSet { // OVERVIEW omitted for space
    int      elts[MAXELTS];
    int      numElts;
    int indexOf(int v); // RME omitted for space
public:
    void insert(int v); void remove(int v); // RME omitted
    bool query(int v); int size();          // RME omitted
};
```

```
int IntSet::indexOf(int v) {
    for (int i = 0; i < numElts; i++) {
        if (elts[i] == v) return i;
    }
    return MAXELTS;
}
```

Abstract Data Types

Using Classes

```
const int MAXELTS = 100;

class IntSet { // OVERVIEW omitted for space
    int      elts[MAXELTS];
    int      numElts;
    int indexOf(int v); // RME omitted for space
public:
    void insert(int v); void remove(int v); // RME omitted
    bool query(int v);  int  size();        // RME omitted
};
```

With `indexOf`, `query` is trivial...

```
bool IntSet::query(int v)    {
    return (indexOf(v) != MAXELTS);
}
```

Abstract Data Types

Using Classes

- The code for `insert` is not much more difficult than query:
 - First look for the `indexOf` the element to insert.
 - If it doesn't exist, we need to add this element to the **end** of the array.
 - What is the index of the current “end” ?



- Place the element in the next slot and **update** `numElements`.
- The only exception to this is if `numElements` already equals `MAXELTS`.

Abstract Data Types

Using Classes

```
const int MAXELTS = 100;
class IntSet { // OVERVIEW omitted for space
    int      elts[MAXELTS];
    int      numElts;
    int indexOf(int v); // RME omitted for space
public:
    void insert(int v);  void remove(int v); // RME omitted
    bool query(int v);  int  size();          // RME omitted
};
```

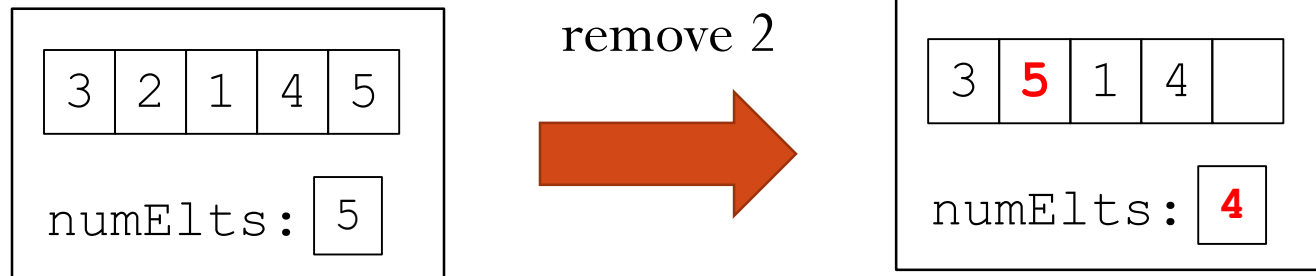
```
void IntSet::insert(int v) {
    if (indexOf(v) == MAXELTS) {
        if (numElts == MAXELTS) throw MAXELTS;
        elts[numElts++] = v;
    }
}
```


How about Remove?

- If the element (called the `victim`) is in the array, we have to remove it leaving a "hole" in the array.
- What representation invariants are violated?
 - How can we fix them?

How about Remove?

- Instead of moving each element after the victim to the left by one position, pick up the current "last" element and move it to the hole.
- This also breaks the invariant on `numElems`, so we must fix it.



Abstract Data Types

Using Classes

```
void IntSet::remove(int v) {  
    int victim = indexOf(v);  
    if (victim != MAXELTS) {  
        elts[victim] = elts[numElts-1];  
        numElts--;  
    }  
}
```

Abstract Data Types

Using Classes

- Question: There is one problem with our implementation. What is it?
- Hint: Consider the newly-created set:

```
IntSet s;
```

What does the computer actually create when we declare `s`?

Abstract Data Types

Using Classes

- Question: There is one problem with our implementation. What is it?
- Answer: On creation, `s`'s data members are **uninitialized**!
- This means that the value of `numElements` could be a random value, but our representational invariant says it must be zero!
- How can we fix this?

Abstract Data Types

Automatically Initializing Classes

- Using **constructor**!
- The constructor (really, the **default** constructor) has the following type signature:

```
class IntSet { // OVERVIEW omitted for space
    ...
    public:
        IntSet();
        // EFFECTS: creates an empty IntSet
    ...
};
```

Abstract Data Types

Automatically Initializing Classes

```
IntSet() ;  
    // EFFECTS: creates an empty IntSet
```

- The name of the function is the same as the name of the class.
- This function doesn't have a return type.
- It also does not take an argument in this case.
- It is guaranteed to be the **first** function called immediately after an object is created.
- It builds a “blank” uninitialized `IntSet` and makes it satisfy the rep invariant.

Abstract Data Types

Automatically Initializing Classes

```
IntSet();  
    // EFFECTS: creates an empty IntSet
```

- Here's how it's written:

```
IntSet::IntSet() : numElts(0)  
{  
}
```


Abstract Data Types

Automatically Initializing Classes

```
IntSet::IntSet()  
    : numElts(0)  
{  
}
```

```
Class_T::Class_T() : anInt(0),  
                    aDouble(1.2),  
                    aString("Yes")  
{  
}
```

- This syntax is called "initialization syntax".
- Each data member is initialized this way.
- **Note**: The order in which elements are initialized is the order they **appear in the definition**, NOT the order in the initialization list. It is a good practice to keep them in the same order to avoid confusion.

Abstract Data Types

Automatically Initializing Classes

- Alternatively, we could write this function as follows, but this is not considered as a good way!

```
IntSet::IntSet()  
{  
    numElts = 0;  
}
```



Not Recommended

Abstract Data Types

A Benefit of Classes

- Now, instead of writing this:

```
void add_one (int a[], int elts);
```

and having to worry about the number of elements in the array. All we have to write is this:

```
void add_one (IntSet& set);
```

and we no longer have to worry about the array and its count being separated.

Const Member Functions

- A slight change to the class definition:

```
const int MAXELTS = 100;
class IntSet {
    int      elts[MAXELTS];
    int      numElts;
    int      indexOf(int v) const;

public:
    void      insert(int v);
    void      remove(int v);
    bool      query(int v) const;
    int      size() const;
};
```

Const Member Functions

```
int size() const;
```

- Each member function of a class has an extra, implicit parameter named **this**.
 - “**this**” is a pointer to the current instance on which the function is invoked.
- **const** keyword modifies the implicit **this** pointer: **this** is now a pointer to a **const instance**.
 - Means: the member function **size()** cannot change the object on which **size()** is called.
 - By its definition, **size()** shouldn't change the object! Adding **const** keyword prevents any accidental change.
 - It is a good practice to add **const** keyword when possible!

Const Member Functions

- Implement **size()**

```
int IntSet::size() const {  
    return numElts;  
}
```

The function body is the same as before.

- A **const** object can only call its **const** member functions!

```
const IntSet is;  
cout << is.size(); ✓  
is.insert(2); ✗
```

Const Member Functions

- If a const member function calls other **member** functions, they must be **const** too!

```
void A::g() const { f(); }
```

```
void A::f() {...}
```



```
void A::f() const {...}
```



Outline

- A Class Example: a Mutable Set of Integers (IntSet)
- Improve the Efficiency of IntSet

Abstract Data Types

Class Exercise

```
int IntSet::indexOf(int v) {  
    for (int i = 0; i < numElts; i++) {  
        if (elts[i] == v) return i;  
    }  
    return MAXELTS;  
}
```

- **Question**: How many elements of the array must `indexOf` examine in the **worst case** if there are 10 elements? If there are 90 elements?

Abstract Data Types

Improving Efficiency

- We say the time for `indexOf` grows **linearly** with the size of the set.
- If there are N elements in the set, we have to examine all N of them **in the worst case**. For large sets that perform lots of queries, this is too expensive!
- Luckily, we can replace this implementation with a different one that can be more efficient. The only change we need to make is to the **representation (implementation)** – the abstraction can stay precisely the same.

Abstract Data Types

Improving Efficiency

- Still use an array to store the elements of the set and the values will still occupy the first `numElements` slots.
- However, now we'll keep the elements in **sorted** order.

Question: What Parts of the Class Should Be Changed?

```
const int MAXELTS = 100;
class IntSet {
    // OVERVIEW: a mutable set of integers
    int elts[MAXELTS];
    int numElts;
    int indexOf(int v) const;
public:
    IntSet();
    void insert(int v);
    void remove(int v);
    bool query(int v) const;
    int size() const;
};
```

Abstract Data Types

Improving Efficiency

- The constructor and size methods don't need to change at all since they just use the `numElts` field.

- `query` also doesn't need to change.

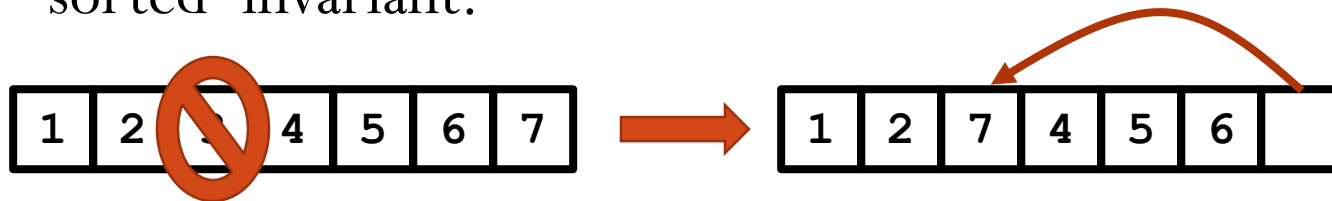
```
bool IntSet::query(int v)    {  
    return (indexOf(v) != MAXELTS);  
}
```

- `indexOf` also doesn't need to change.
- However, `insert` and `remove` do need to change.

Abstract Data Types

Improving Efficiency

- We'll start with the easiest one: `remove`.
- Recall the old version that moved the last element from the end to somewhere in the middle, this will break the new “sorted” invariant.



- Instead of doing a swap, we have to "squish" the array together to cover up the hole.



Abstract Data Types

Improving Efficiency

- How are we going to do the “squish”?
 - Move the element next to the hole to the left leaving a new hole.
 - Keep moving elements until the hole is “off the end” of the elements.

1	2	4	6	7	
---	---	---	---	---	--

remove 4

1	2		6	7	
---	---	--	---	---	--

1	2	6		7	
---	---	---	--	---	--

1	2	6	7		
---	---	---	---	--	--

- We'll reuse the variable `victim` as a loop variable.
- `victim`'s invariant is that it always points at the hole in the array.

Abstract Data Types

Improving Efficiency

```
void IntSet::remove(int v) {  
    int victim = indexOf(v);  
    if (victim != MAXELTS) {  
        // victim points at hole  
        numElts--; // one less element  
        while (victim < numElts) {  
            // ..hole still in the array  
            elts[victim] = elts[victim+1];  
            victim++;  
        }  
    }  
}
```


Abstract Data Types

Improving Efficiency

- We also have to change `insert` since it currently just places the new element at the end of the array. This also will break the new “sorted” invariant.



Abstract Data Types

Improving Efficiency

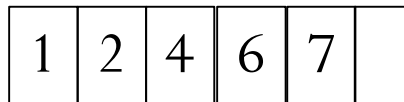
- How are we going to do the insert?
 - Start by moving the last element to the right by one position.
 - Repeat this process until the correct location is found to insert the new element.
 - Stop if the start of the array is reached or the element is sorted.
 - We'll need a new loop variable called `candidate` to track this movement.
 - It's invariant in that it always points to the next element that might have to move to the right.

Abstract Data Types

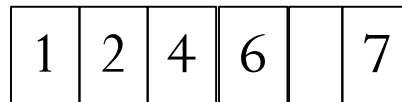
Improving Efficiency

```
void IntSet::insert(int v) {  
    if (indexOf(v) == MAXELTS) { // duplicate not found  
        if (numElts == MAXELTS) throw MAXELTS; // no room  
        int cand = numElts-1; // last element  
        while ((cand >= 0) && elts[cand] > v) {  
            elts[cand+1] = elts[cand];  
            cand--;  
        }  
        // Now, cand points to the left of the "gap".  
        elts[cand+1] = v;  
        numElts++; // repair invariant  
    }  
}
```

insert 5



↑
cand



↑
cand



↑
cand

Abstract Data Types

Improving Efficiency

```
void IntSet::insert(int v) {  
    if (indexOf(v) == MAXELTS) { // duplicate not found  
        if (numElts == MAXELTS) throw MAXELTS; // no room  
        int cand = numElts-1; // last element  
        while ((cand >= 0) && elts[cand] > v) {  
            elts[cand+1] = elts[cand];  
            cand--;  
        }  
        // Now, cand points to the  
        elts[cand+1] = v;  
        numElts++; // repair invariant  
    }  
}
```

Note: We are using the "short-circuit" property of &&. If cand is not greater than or equal to zero, we never evaluate the right-hand clause.

Abstract Data Types

Improving Efficiency

```
void IntSet::insert(int v) {  
    if (indexOf(v) == MAXELTS) { // duplicate not found  
        if (numElts == MAXELTS) throw MAXELTS; // no room  
        int cand = numElts-1; // largest (last) element  
        while ((cand >= 0) && elts[cand] > v) {  
            elts[cand+1] = elts[cand];  
            cand--;  
        }  
        // Now, cand points to the left of the "gap".  
        elts[cand+1] = v;  
        numElts++; // repair invariant  
    }  
}
```

Question: What is the situation when the loop terminates due to `cand < 0`? Is our implementation correct?

Abstract Data Types

Improving Efficiency

- **Question**: Do we have to change `indexOf`?

```
int IntSet::indexOf(int v) {  
    for (int i = 0; i < numElts; i++) {  
        if (elts[i] == v) return i;  
    }  
    return MAXELTS;  
}
```

Abstract Data Types

Improving Efficiency

- **Question**: Do we have to change `indexOf`?
- **Answer**: No, but it can be made more efficient with the new representation.
- **How?** Using **binary search**! (The array is sorted)

```
int IntSet::indexOf(int v) {  
    for (int i = 0; i < numElts; i++) {  
        if (elts[i] == v) return i;  
    }  
    return MAXELTS;  
}
```

Abstract Data Types

Complexity

	<u>Unsorted</u>	<u>Sorted</u>
query	$O(N)$?
insert	?	?
remove	?	?

Abstract Data Types

Complexity

	<u>Unsorted</u>	<u>Sorted</u>
query	$O(N)$	$O(\log N)$
insert	$O(N)$	$O(N)$
remove	$O(N)$	$O(N)$

insert and remove are still **linear**, because they may have to "swap" an element to the beginning/end of the array.

Abstract Data Types

Complexity

	<u>Unsorted</u>	<u>Sorted</u>
query	$O(N)$	$O(\log N)$
insert	$O(N)$	$O(N)$
remove	$O(N)$	$O(N)$

- If you are going to do more searching than inserting/removing, you should use the "sorted array" version, because `query` is faster there.
- However, if `query` is relatively rare, you may as well use the "unsorted" version. It's "about the same as" the sorted version for `insert` and `remove`, but it's MUCH simpler!

References

- Constructor
 - Problem Solving with C++, 8th Edition, Chapter 10.2
- const Member Function
 - C++ Primer, 4th Edition, Chapter 7.7.1