

# Ve 280

Programming and Introductory Data Structures

Improve ADT Efficiency;  
Midterm Review

# Abstract Data Types

## Improving Efficiency

```
void IntSet::insert(int v) {  
    if (indexOf(v) == MAXELTS) { // duplicate not found  
        if (numElts == MAXELTS) throw MAXELTS; // no room  
        int cand = numElts-1; // largest (last) element  
        while ((cand >= 0) && elts[cand] > v) {  
            elts[cand+1] = elts[cand];  
            cand--;  
        }  
        // Now, cand points to the left of the "gap".  
        elts[cand+1] = v;  
        numElts++; // repair invariant  
    }  
}
```

Question: What is the situation when the loop terminates due to `cand < 0`? Is our implementation correct?

# Abstract Data Types

## Improving Efficiency

- Question: Do we have to change `indexOf`?

```
int IntSet::indexOf(int v) {  
    for (int i = 0; i < numElts; i++) {  
        if (elts[i] == v) return i;  
    }  
    return MAXELTS;  
}
```

# Abstract Data Types

## Improving Efficiency

- **Question**: Do we have to change `indexOf`?
- **Answer**: No, but it can be made more efficient with the new representation.
- **How?** Using **binary search**! (The array is sorted)

```
int IntSet::indexOf(int v) {  
    for (int i = 0; i < numElts; i++) {  
        if (elts[i] == v) return i;  
    }  
    return MAXELTS;  
}
```

# Abstract Data Types

Complexity

	<u>Unsorted</u>	<u>Sorted</u>
query	$O(N)$	?
insert	?	?
remove	?	?

# Abstract Data Types

## Complexity

	<u>Unsorted</u>	<u>Sorted</u>
query	$O(N)$	$O(\log N)$
insert	$O(N)$	$O(N)$
remove	$O(N)$	$O(N)$

insert and remove are still **linear**, because they may have to "swap" an element to the beginning/end of the array.

# Abstract Data Types

## Complexity

	<u>Unsorted</u>	<u>Sorted</u>
query	$O(N)$	$O(\log N)$
insert	$O(N)$	$O(N)$
remove	$O(N)$	$O(N)$

- If you are going to do more searching than inserting/removing, you should use the "sorted array" version, because `query` is faster there.
- However, if `query` is relatively rare, you may as well use the "unsorted" version. It's "about the same as" the sorted version for `insert` and `remove`, but it's MUCH simpler!

# Midterm Review

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# Midterm

- 10:00 am – 11:40 am, June 27<sup>th</sup>, 2017
- Find your seat on Canvas
- Closed book and closed notes
- No electronic devices are allowed
  - These include laptops and cell phones
  - We will show a clock on the screen

# Midterm

- Written exam
  - A number of questions which only require you to provide a very short answer
  - A few questions which require you to write code on the paper. Be clear!
- Abide by the **Honor Code!**

# Midterm Topics

- Linux Commands
- Compiling and Developing Program on Linux
- C++ Basics: Pointers, References, const Qualifier
- Procedural Abstraction and Specification Comments
- Recursion
- Function Pointers
- enum Type
- Program Taking Arguments
- I/O Streams
- Testing/Debugging
- Exception
- Class Basics

Lecture 1 to this lecture

# Linux Commands

- `cd; ls; mkdir; rmdir;`
- `cp; mv; rm;`
- `nano; gedit;`
- `less;`
- `diff; man; ...`
- I/O redirection
  - `<, >`
- Command options
  - `ls -l; cp -r dir1 dir2; ...`
- Wildcard: `*`
  - `cp *.h dir/`

# Compiling Program on Linux

- Write the source code, for example, using **gedit**
- Compile the program: **g++ -o program source.cpp**
- Run the program: **./program**
- Compile multiple source files:
  - **g++ -o program src1.cpp src2.cpp src3.cpp**
  - E.g., **g++ -o run\_add run\_add.cpp add.cpp**
- Header guard: avoiding multiple inclusions

```
// add.h
#ifndef ADD_H
#define ADD_H
int add(int a, int b);
#endif
```

- What happens if the .h file is included **first** time?
- What happens if the .h file is included **second** time?

# A Better Way of Compiling: Makefile

`all: run_add`

- The file name is “**Makefile**”
- Type “**make**” on command-line

`run_add: run_add.o add.o`

`g++ -o run_add run_add.o add.o`

`run_add.o: run_add.cpp`

`g++ -c run_add.cpp`

`add.o: add.cpp`

`g++ -c add.cpp`

`clean:`

`rm -f run_add *.o`

A Rule

Target: Dependency  
<Tab> Command



Don't forget the Tab!

Dependency: A list of files  
that the target depends on

# Function Call Mechanisms

There are two function call mechanisms:

1. Call-by-value
2. Call-by-reference

What will a be?

```
void f(int x)
{
    x *= 2;
}
```

```
int main()
{
    ...
    int a=4;
    f(a);
    ...
}
```

```
void f(int& x)
{
    x *= 2;
}
```

```
int main()
{
    ...
    int a=4;
    f(a);
    ...
}
```

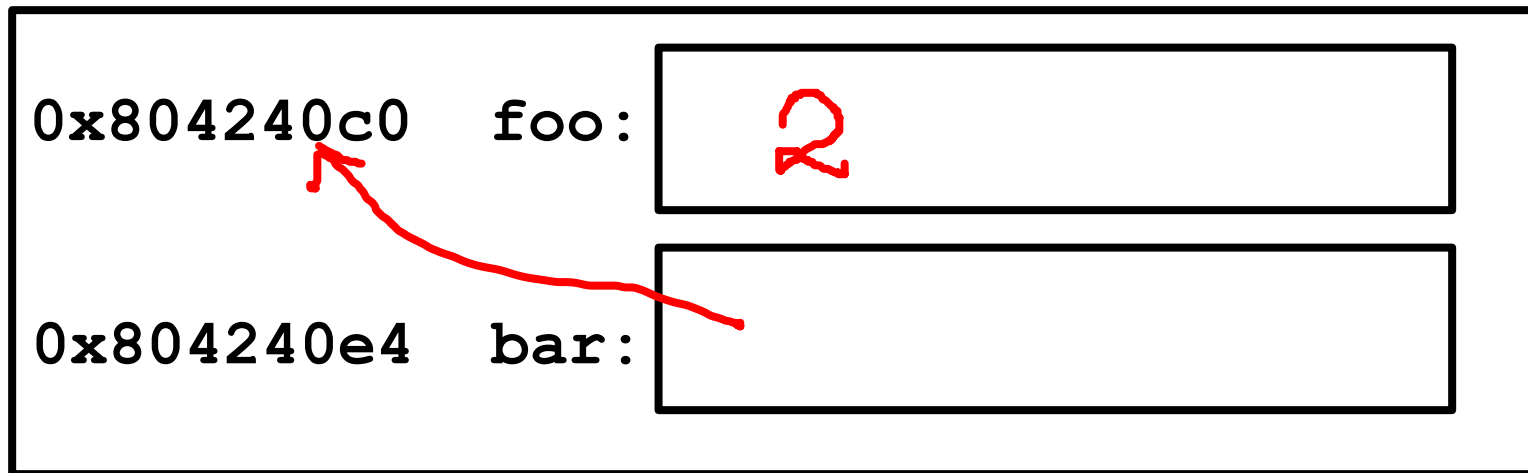
# Pointers

```
int foo = 1;
```

```
int *bar;
```

```
bar = &foo; // addressing operation
```

```
*bar = 2; // dereference operation
```





# References

- An alternative name for an object

```
int iVal = 1024;  
int &refVal = iVal;
```

- Reference **must be initialized** using a **variable** of the same type.

# References Versus Pointers

## Example

```
int x = 0;  
int &r = x;  
int y = 1;  
r = y;  
r = 2;
```

What's the final values of  
x, y, and r?

x = 2, y = 1, r = 2

```
int x = 0;  
int *p = &x;  
int y = 1;  
p = &y;  
*p = 2;
```

What's the final values of  
x, y, and \*p?

x = 0, y = 2, \*p = 2

# const Qualifier

- Once you defined a constant variable, it cannot be modified later on.

- `const int a = 10;`  
`a = 11; // Error`

- Because we cannot subsequently change the value of an object declared to be const, we must initialize it when it is defined:

- `const int i;`  
`// Error: i is an uninitialized const`

# const Reference

```
int avg_exam(const struct Grades & gr) {  
    return (gr.midterm+gr.final)/2;  
}
```

- It gives us the best of both worlds:
  - We don't have the expense of a copy.
  - We have the safety guarantee that the function cannot change the caller's state. Compiler will catch the error of accident change!

# const Pointers

- When you have pointers, there are two things you might change:
  1. The value of the pointer.
  2. The value of the object to which the pointer points.
- Either (or both) can be made unchangeable:

```
const T *p;    // "T" (the pointed-to object)
               // pointer to const // cannot be changed by pointer p
T *const p;    // "p" (the pointer) cannot be
               // const pointer  // changed
const T *const p; // neither can be changed.
```

# Pointers to const

## Example

```
int a = 53;
const int *cptr = &a;
    // OK: A pointer to a const object
    // can be assigned the address of a
    // nonconst object
*cptr = 42;
    // ERROR: We cannot use a pointer to
    // const to change the underlying
    // object.
a = 28 // OK
int b = 39;
cptr = &b; // OK: the value in the pointer
           // can be changed.
```

# const Pointers

## Example


```
int a = 53;
int *const cptr = &a;
    // OK: initialization
*cptr = 42;
    // OK: We can use a const pointer to
    // change the underlying object.
int b = 39;
cptr = &b;
    // ERROR: We cannot change the value of
    // a const pointer.
```

# Pointer to const versus Normal Pointer

- Pointers-to-const-T are **not the same** type as pointers-to-T.
- You can use a pointer-to-T anywhere you expect a pointer-to-const-T, but NOT vice versa.


```
int const_ptr(const int *ptr)
{
    ...
}

int main()
{
    int a = 0;
    int *b = &a;
    const_ptr(b);
}
```



```
int nonconst_ptr(int *ptr)
{
    ...
}

int main()
{
    int a = 0;
    const int *b = &a;
    nonconst_ptr(b);
}
```





# Abstraction

- Abstraction
  - Provides only those details that matter.
  - Eliminates unnecessary details and reduces complexity.
- Example: Multiplication algorithm
  - Many ways to do: table lookup, summing, etc.
  - Each looks quite different, but they do the **same** thing.
  - In general, a user won't care how it's done, just that it multiplies.

# Procedural Abstraction

- Two important properties of procedural abstraction
  - **Local**: the implementation of an abstraction does not depend on any other abstraction implementation.
  - **Substitutable**: you can replace one (correct) implementation of an abstraction with another (correct) one, and no callers of that abstraction will need to be modified.

# Procedural Abstraction

## Specification Comments

- We describe procedural abstraction by specification comments.
- There are three clauses of specification comments:
  - **REQUIRES**: the pre-conditions that must hold, if any.
  - **MODIFIES**: how inputs are modified, if any.
  - **EFFECTS**: what the procedure computes given legal inputs.
- Note that the first two clauses have an “**if any**”, which means they may be empty, in which case you may omit them.

# Call Stacks

How a function call really works

- When a function is called, an activation record (also known as stack frame) is created. It holds the function's **formal parameters** and **local variables**.
- The **activation record** for **the current** invocation is added to the “top” of the stack.
- When that function returns, its **activation record** is removed from the “top” of the stack.



double add(double a, double b): a = 1, b = 0, result = 0

double sin(double x): x = 1, result = 0

int main(): x = 1, sinResult = 0

# Recursion

$$n! = \begin{cases} 1 & (n == 0) \\ n * (n-1)! & (n > 0) \end{cases}$$

---

```
int factorial (int n)
    // REQUIRES: n >= 0
    // EFFECTS:  computes n!

1.  {
2.      if (n == 0) {
3.          return 1;  // 'base case'
4.      } else {
5.          return n*factorial(n-1); // 'recursive step'
6.      }
    }
```

# Recursion

Writing a function for the general case

- Treat it like an inductive proof.
- To write a correct recursive function, do two things:
  1. Identify the “trivial” case (or cases), and write them explicitly.
  2. For all other cases, first assume there is a function that can solve smaller versions of the same problem, then figure out how to get from the smaller solution to the bigger one.

# Recursive Helper Function

- Sometimes it is easier to find a recursive solution to a problem if you change the original problem slightly, and then solve that problem using a **recursive helper function**.

```
soln()  
{  
    ...  
    soln_helper();  
    ...  
}
```

```
soln_helper()  
{  
    ...  
    soln_helper();  
    ...  
}
```

# Function Pointers

## Motivation

- If you were asked to write a function to add all the elements in a list, and another to multiply all the elements in a list, your functions would be almost exactly **the same**.
- Writing almost the exact same function twice is almost certainly a bad idea

**Function pointers to the rescue!**



# Function Pointers

## A first look

```
int min(int a, int b);  
    // EFFECTS: returns the smaller of a and b.  
int max(int a, int b);  
    // EFFECTS: returns the larger of a and b.
```

- These two functions have precisely the same type signature:
  - They both take two integers, and return an integer.
- Of course, they do completely different things:
  - One returns a min and one returns a max.
  - **However, from a syntactic point of view, you call either of them the same way.**

# Function Pointers

## Basic Format

- Declaration

```
int    (*foo) (int, int);
```

- Once defined, we can assign it to a function that has **the same type signature**

```
int min(int a, int b);  
foo = min;
```

- Furthermore, after assigning min to foo, we can just call it as follows:

```
foo(3, 5)
```

...and we'll get back 3!

# Enum Type

- Define an enumeration type as follows:

```
enum Suit_t {CLUBS, DIAMONDS,  
             HEARTS, SPADES};
```

- Define variables of this enum type:

```
enum Suit_t suit;
```

- You can initialize them as:

```
enum Suit_t suit = DIAMONDS;
```

- Once you have such an enum type defined, you can use it as an argument for a function.

# Enum Type

- If you write

```
enum Suit_t {CLUBS, DIAMONDS,  
             HEARTS, SPADES};
```

then numerically

```
CLUBS = 0, DIAMONDS = 1,  
HEARTS = 2, SPADES = 3
```

- Using this fact, it will sometimes make life easier

```
enum Suit_t s = CLUBS;  
const string suitname[] = {"clubs",  
                           "diamonds", "hearts", "spades"};  
cout << "suit s is " << suitname[s];
```

# Passing Arguments to a Program

- Programs can take arguments.

**diff file1 file2**

- Arguments are passed to the program through `main()` function.
- We need to change the argument list of `main()`:
  - `int main(int argc, char *argv[])`
- `argv` stores the array of C-strings that user inputs.
  - `argv[0]` is the name of the program being executed.
- `argc` is the number of strings in the array

# I/O Streams

- Output Stream `cout`
  - Insertion operator `<<`
- Input Stream `cin`
  - extraction operator `>>`
  - `getline(cin, str)`
  - `cin.get(ch)`
  - Failed input stream: check stream state `if (cin)`
- `cout` and `cin` streams are buffered.

# I/O Streams

- File Stream
  - `ifstream; ofstream`
  - Opening a file: `ifstream.open("myText.txt");`
  - extraction `>>` ; insertion `<<`
- String Stream
  - `istringstream; ostringstream`
  - extraction `>>` ; insertion `<<`
  - Assign a string to an input string stream  
`istream.str(line);`
  - fetch the string value from an output string stream  
`ostream.str();`

# Testing

- Be skeptical!
- Incremental testing
- Five Steps:
  1. Understand the specification
  2. Identify the required behaviors
  3. Write specific tests
    - **Simple inputs**
    - **Boundary conditions**
    - **Nonsense**
  4. Know the answers in advance
  5. Include stress tests



# Debugging Using Assert

- Using the `assert` function
  - The `assert` function is a special function, which takes a Boolean argument.
  - If the argument is **true**, `assert()` does nothing.
  - If the argument is **false**, `assert()` causes your program to stop, printing an **error message** to the `cerr` stream.
- `assert` for the condition that should hold.

# Exceptions

- Exceptions and exception handling mechanism



- **Exception propagation** mechanism: where to find the handler

# Exceptions

## Exception Handling

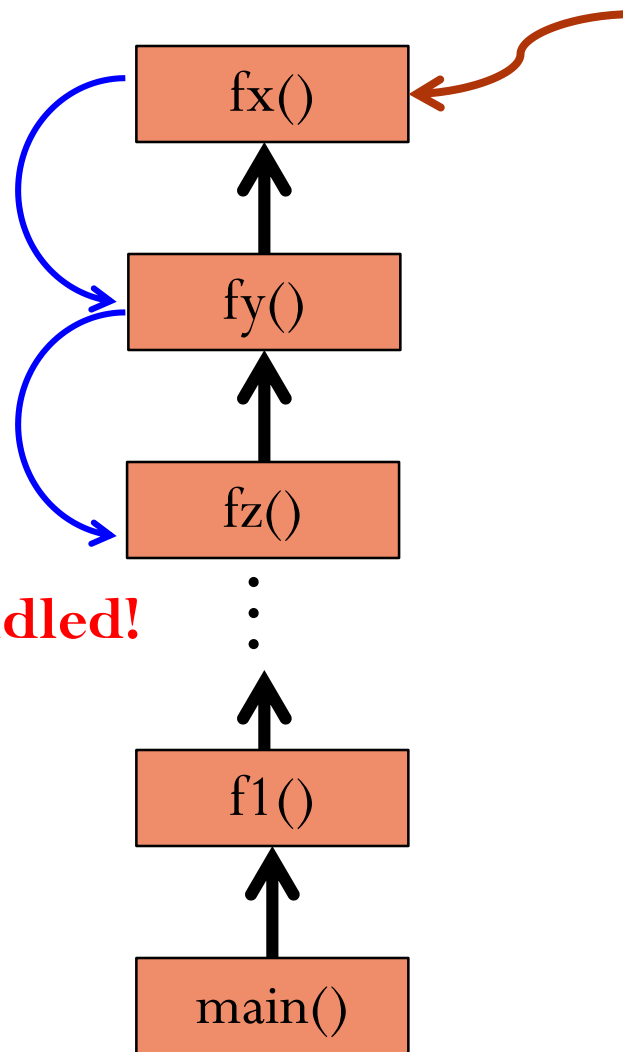
handler in fx()? **No!**

handler in fy()? **No!**

handler in fz()? **Yes!**

**Exception handled!**

**exception  
occurs**



# Exceptions

- Throwing an exception
- Catching an exception
- Exceptions have **types** and **objects**.
  - **throw** **errorObj**;
- Exceptions Handling in C++

```
void foo() {  
    try {  
        catch (Type var) {  
    }  
}
```

# Abstract Data Types

- The role of a type:
  - The set of values that can be represented by items of the type
  - The set of operations that can be performed on items of the type.
- An abstract data type provides an **abstract description** of **values** and **operations**.
- Advantages: Information hiding and encapsulation.

# C++ Classes

- Data members and function members are defined in a single entity.
- **Public** versus **private** members.
- Defining a class type.
- Class object as a function argument: pass by value

# C++ Classes

- **Constructor** for initialization: `IntSet () ;`

- Initialization syntax:

```
IntSet::IntSet () : numElts (0)
{ }
```

- const member function: `int size() const;`
  - Means: the member function **size()** cannot change the object on which **size()** is called.
  - Syntax: if a const member function calls other **member** functions, they must be **const** too!

```
void A::g() const { f(); }
```

```
void A::f() { ... }
```



```
void A::f() const { ... }
```

