

CRYPTOGRAPHY  
ACADEMIC YEAR 2025-2026  
HOMEWORK I  
OCTOBER 10TH, 2025

Please notice that:

- Exercises are meant to be solved *individually*.
- Solutions should be typeset in  $\text{\LaTeX}$ , and uploaded, in pdf format, to <http://virtuale.unibo.it>. Students are encouraged to use the template `Homework-template-2526.tex`, which can be retrieved from <http://virtuale.unibo.it> itself.
- The deadline for uploading the solutions is Monday, October 17th, at midnight CET.

In all the exercises below, if  $y \in \{0, 1\}^m$  and  $n \leq m$ , then  $y|_n$  is the binary string obtained by considering the first  $n$  bits of  $y$ .

**Exercise 1.**

The *scytale* (see, <https://en.wikipedia.org/wiki/Scytale>) is a classical cipher from ancient Greece. Define it as a triple of algorithms  $\Pi = (\text{Gen}, \text{Enc}, \text{Dec})$  as we did in Section 2, and prove that  $\Pi$  is not perfectly secure.

**Exercise 2.**

Consider a pseudorandom generator  $G$  with expansion factor  $\ell$ , and let  $\ell'$  be any polynomial such that  $\ell(n) > \ell'(n) > n$ . Consider the function  $H$  defined as  $H(x) = G(x)|_{\ell'(|x|)}$ . Prove that  $H$  is still a pseudorandom generator.

**Exercise 3.**

Consider the following functions, and identify which ones of them are pseudorandom generators, proving your claims:

$$G_1(x) = x \cdot 010 \quad G_2(x) = F(x, 0^{|x|}) \quad G_3(x) = \begin{cases} x \cdot x & \text{if } |x| \leq 2 \\ F(x, 1^{|x|}) \cdot F(x, 0^{|x|}) & \text{otherwise} \end{cases}$$

Here,  $\cdot$  is string concatenation, and  $F$  is a pseudorandom function. Moreover, prove that none of the following binary functions is a pseudorandom function:

$$F_1(k, x) = x \oplus k \quad F_2(k, m) = G(k)|_{|k|} \oplus m$$

where  $G$  is a pseudorandom generator.