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CSE 565: Software Verification and Validation
Project 2 Report
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Part 1

1. Description of the tool used and the types of coverage it provides

Tool name: Jacoco

Jacoco is a tool. Its main competitive edge is simple and light, compared to other tools on the market. It also provides code coverage metrics for statement, decision, class and method. Furthermore, it can also be used for code delta testing (if implemented in IntelliJ). Last but not least, it has 2 additional features: output reports and pinpointing uncovered lines.

2. Set of test cases:

#	Item	Input
1	candy	21
2	candy	20
3	candy	19
4	coke	24
5	coffee	44

Though the requirements ask for 100% decision coverage, in reality it only covers at most 93% of all the branches. The achievement of 15/16 decision coverage can be shown in the snapshot. This project creates the control flow (appendix 2.A, 2.B) to construct the data flows. By systemically mapping the data flows in the appendix 3 to the test cases we obtain the data flow coverage mapping table (appendix 1). Using the aforementioned table(appendix 1), we pinpointed the predicate PVI whose False result is unreachable. We will now prove that achieving the last one of the 16 decisions is impossible.

Proof by contradiction:

Assume the path (PIV: False, PV: False, PVI: False) is reachable.

Let $\text{cost}(\text{item})$ be the cost of an item of 3 item types.

We notice 2 facts:

1. The path requires the decisions of predicates PIV, PV, PVI to be false-false-false regardless of all of the decisions of the other predicates (appendix 1, in red).
2. All of the items cost less than or equal to 45, $\text{cost}(\text{item}) \leq 45$ (appendix 4).

From the path construction in the appendices 2,3, we have the following:

def1(input1) -> PIV(input1) -> PV(input1) -> PVI(input1)

For these to be false, all of the following conditions must be true:

(1) $\text{input1} \leq \text{cost}(\text{item})$

(2) $\text{input1} \neq \text{cost}(\text{item})$

(3) $\text{input1} \geq 45$

Combining (1), (2) and (3), we have the inequality:













$45 \leq \text{input1} < \text{cost}(\text{item}) \Leftrightarrow 45 < \text{cost}(\text{item})$

The premise $45 < \text{cost}(\text{item})$ contradicts the fact that all of the items cost less than 45 (which is that $45 \geq \text{cost}(\text{item})$).

Therefore, path (PIV: False, PV: False, PVI: False) is unreachable

Note that: this was done in accordance with the requirement in the CSE 565 Project 2_Analyzing Code Coverage.pdf file. Since the clarification statement during the live event was made late and the result of the proof above did not affect the overall quality of the project, this project decided not to make changes to the project report.

3. Screenshot showing the coverage achieved for the test cases developed

Element	Class, % ▾	Method, %	Line, %	Branch, %
 VendingMachine	100% (1/1)	100% (2/2)	100% (32/32)	93% (15/16)
 com				
 images				
 java				
 javax				
 jdk				
 kotlin				
 META-INF				
 netscape				
 org				
 sun				
 toolbarButtonGraphics				

4. Evaluation of the tool's usefulness

With the tool, software testers can reduce the workload by automating the code coverage results of the test cases. It also supports proof of correctness of the actual results of the test cases against the expected results which were created during the planning of the data flow coverage and control flow coverage.

Part 2:

1. Description of the tool used and the types of analysis it provides

Tool name: IntelliJ IDEA code inspection

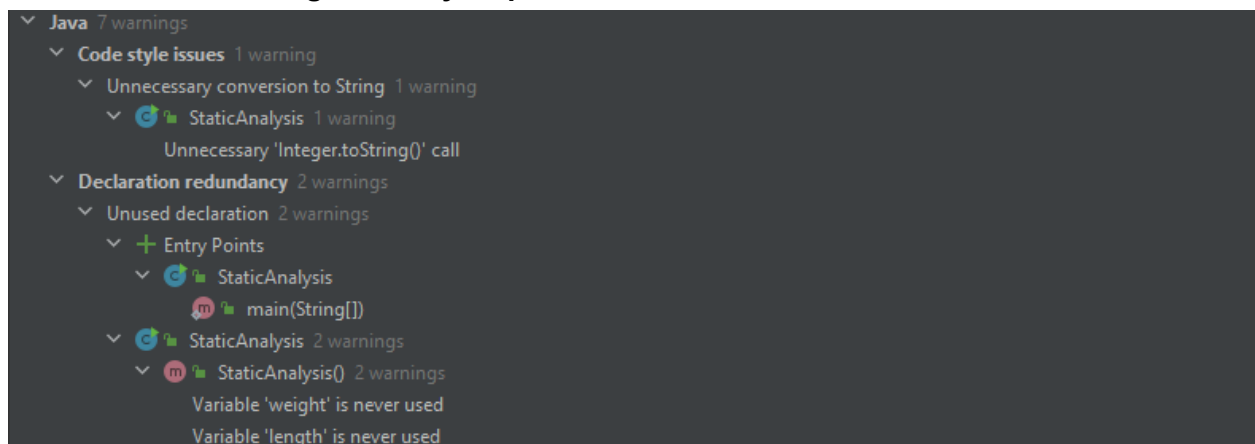
IntelliJ IDEA code inspection, a feature embedded into IntelliJ IDEA, is a static code analyzer. It provides major analyses such as: data flow anomaly, code analysis, etc.

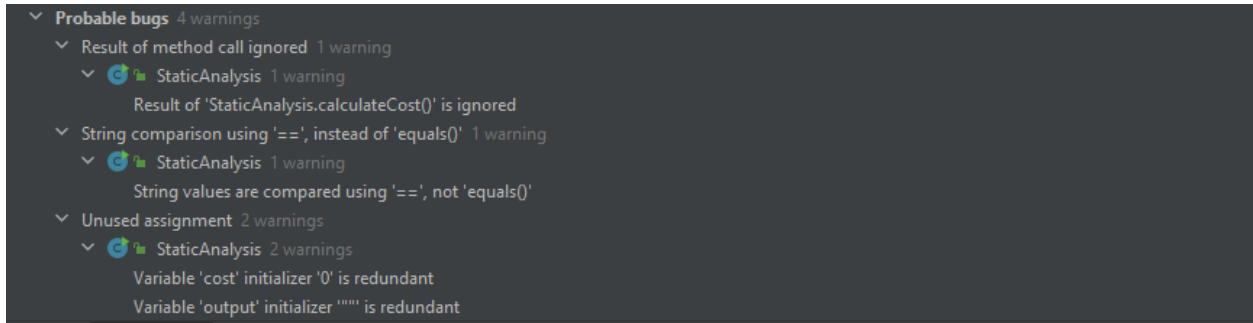
2. Description of the two data flow anomalies

Line	Code	Description
8	<code>calculateCost(5, 10, "Electronics");</code>	The variable output is returned without a container or being used.
15	<code>int weight = 0;</code>	Variable weight is declared and defined without being used
16	<code>String length = "";</code>	Variable length is declared and defined without being used
20	<code>int cost = 0;</code>	the value 0 of cost is never used
21	<code>String output = "";</code>	the value "" of output is never used

Note: there are other warnings but not related to the data flow anomalies, and this project has captured more than 2 anomalies required.

3. Screenshot showing the analysis performed





4. Evaluation of the tool's usefulness

It helps software testers and programmers to prevent the errors from lurking in the program. It reduces time, cost and resources to find the bugs via test cases.

Appendix

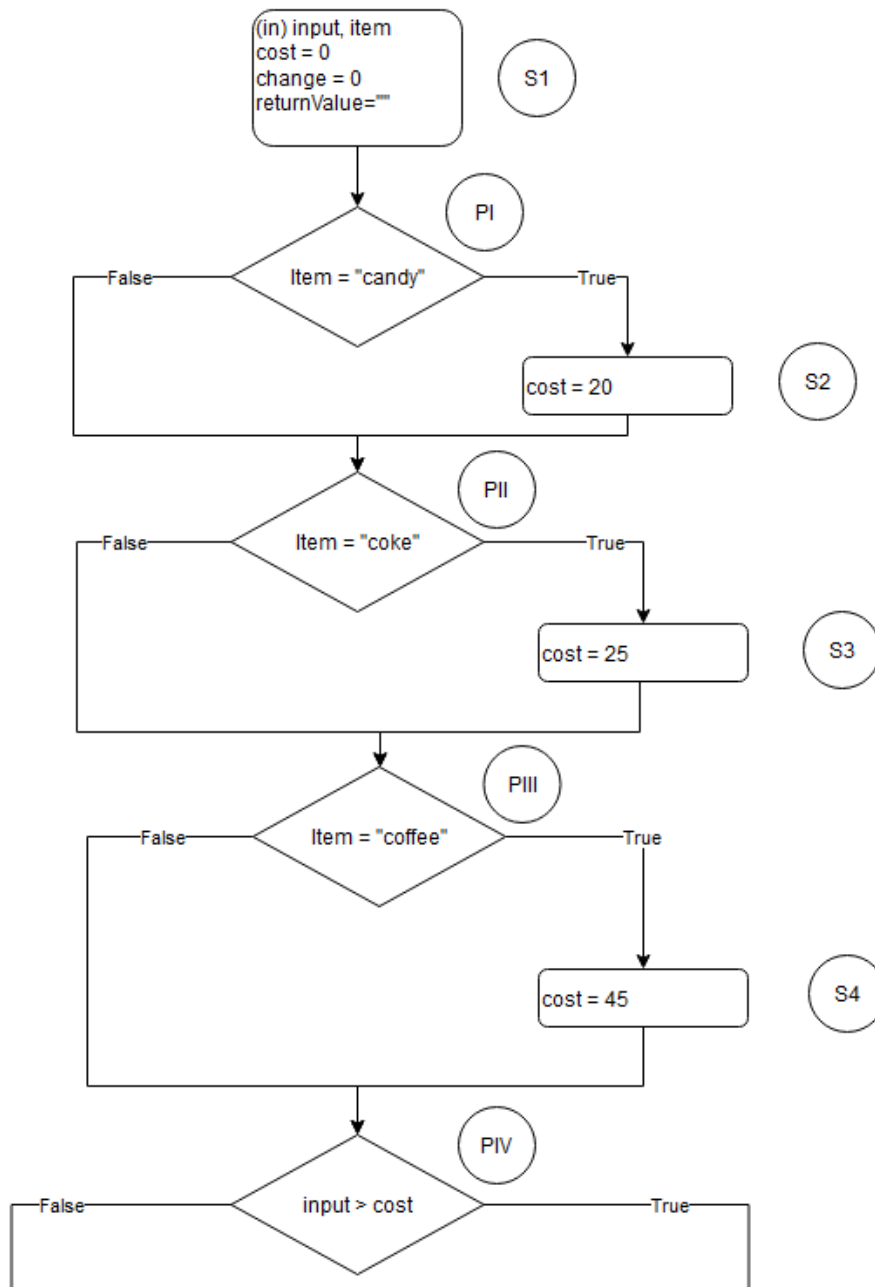
Appendix 1: Data flow path mapping table

Test case	1	2	3	4	5
def1input->PIV	T	F	F	F	F
def1input->PV	NA	T	F	F	F
def1input->C5	x				
def1input->C6			x	x	x
def1input->PVI			T	T	T
def1input->PVII			T	T	F
def1input->PVIII			T	F	F
def1(item) ->PI	T	T	T	F	F
def1(item) ->PII	F	F	F	T	F
def1(item) ->PIII	F	F		F	T
def2(cost) ->PIV	T	F		F	F
def2(cost) ->PV	NA	T	F	F	F
def2(cost) ->C5	x				
def2(cost) ->C6			x		
def3(cost) -> PIV				F	
def3(cost) -> PV				F	
def3(cost) -> C5					
def3(cost) -> C6				x	
def4(cost) -> PIV					F
def4(cost) -> PV					F
def4(cost) -> C5					
def4(cost) -> C6					x
def5(change)->C5	x				
def5(returnValue)->C11	x				

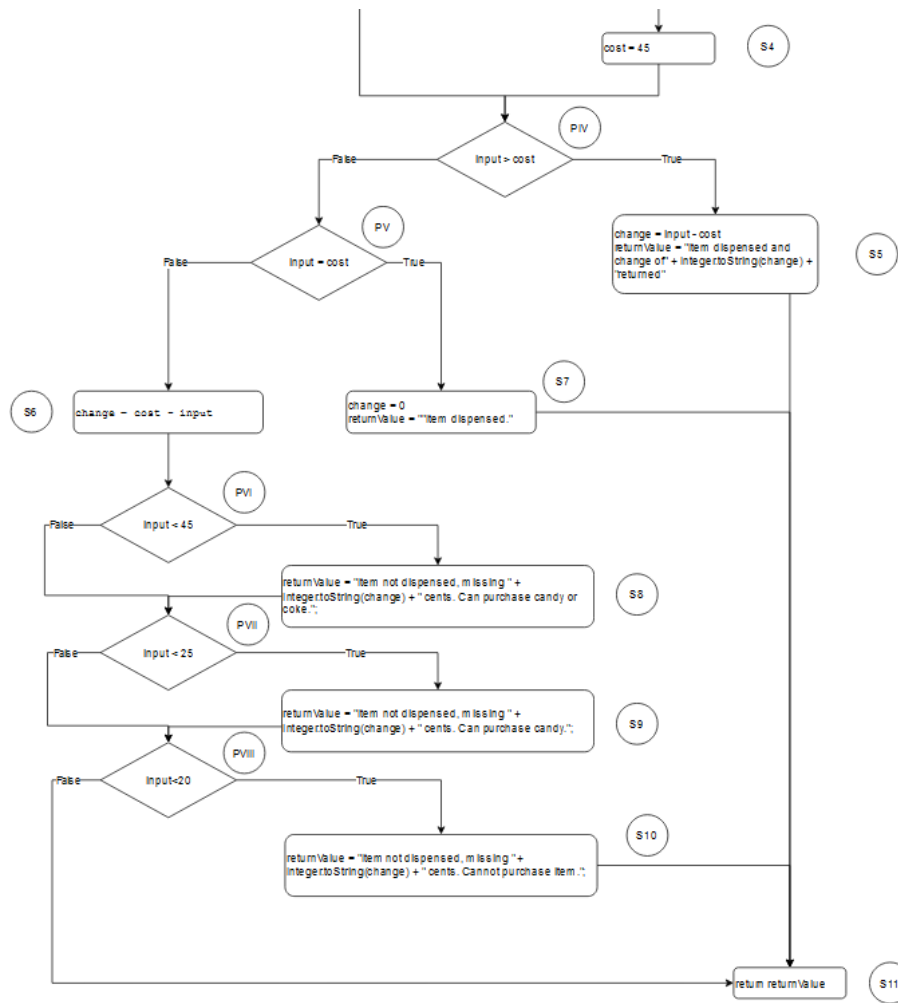
def6(change) -> C8			x	x	x
def6(change) -> C9			x	x	
def6(change) -> C10			x		
def7(returnValue) -> C11		x			
def8(returnValue) -> C11					x
def9(returnValue) -> C11				x	
def10(returnValue) -> C11			x		
def1(change)	N/A	N/A	N/A	N/A	N/A
def1(cost)	N/A	N/A	N/A	N/A	N/A
def1(returnValue)	N/A	N/A	N/A	N/A	N/A
def7(change)	N/A	N/A	N/A	N/A	N/A

Appendix 2

Appendix 2.A: Control flow diagram (Part A)



Appendix 2.B: Control Flow Diagram (Part B)



Appendix 3: Data flow paths

def1(input1) -> PIV(input1) | PV(input1) | C5(input1) | C6(input1) | PVI(input1) | PVII(input1) | PVIII(input1)

def1(item1) -> PI(item1) | PII(item1) | PIII(item1)

def1(change1)

def1(cost1)

def1(returnValue1)

def2(cost2) -> PIV(cost2) | PV(cost2) | C5(cost2) | C6(cost2)

def3(cost3) -> PIV(cost3) | PV(cost3) | C5(cost3) | C6(cost3)

def4(cost4) -> PIV(cost4) | PV(cost4) | C5(cost4) | C6(cost4)

def5(change5) -> C5(change5)

def5(returnValue5) -> C11(returnValue5)

def6(change) -> C8(change6) | C9(change6) | C10(change6)

def7(change7) ->

def7(returnValue7) -> C11(returnValue7)

def8(returnValue8) -> C11(returnValue8)

def9(returnValue9) -> C11(returnValue9)

def10(returnValue10) -> C11(returnValue10)

Appendix 4: Explanation of observation 1 and the premise gained from the symbolic execution

From the result of the symbolic execution of the program(below), we have the maximum cost of all the items at PVI is $\max(\text{cost2}, \text{cost3}, \text{cost4}) = \max(20, 25, 45) = 45$, which means that $(\text{cost}(\text{item}) \leq 45)$.

S1:

input1, item1

cost1 = 0

change1 = 0

returnValue1 = 0

S2

cost2 = 20

S3

cost3 = 25

S4

cost4 = 45

S5

change5 = input1 - cost(2|3|4)

returnValue5 = change5 = input - cost(2|3|4)

S6

change6 = cost1 - input1

S7

change7

returnValue7

S8

returnValue8 = change6 = cost1 - input1

S9

returnValue9 = change6 = cost1 - input1

S10

returnValue10 = change6 = cost1 - input1

S11

returnValue11 = returnValue(5|7|8|9|10)