

Advanced Artificial Intelligence

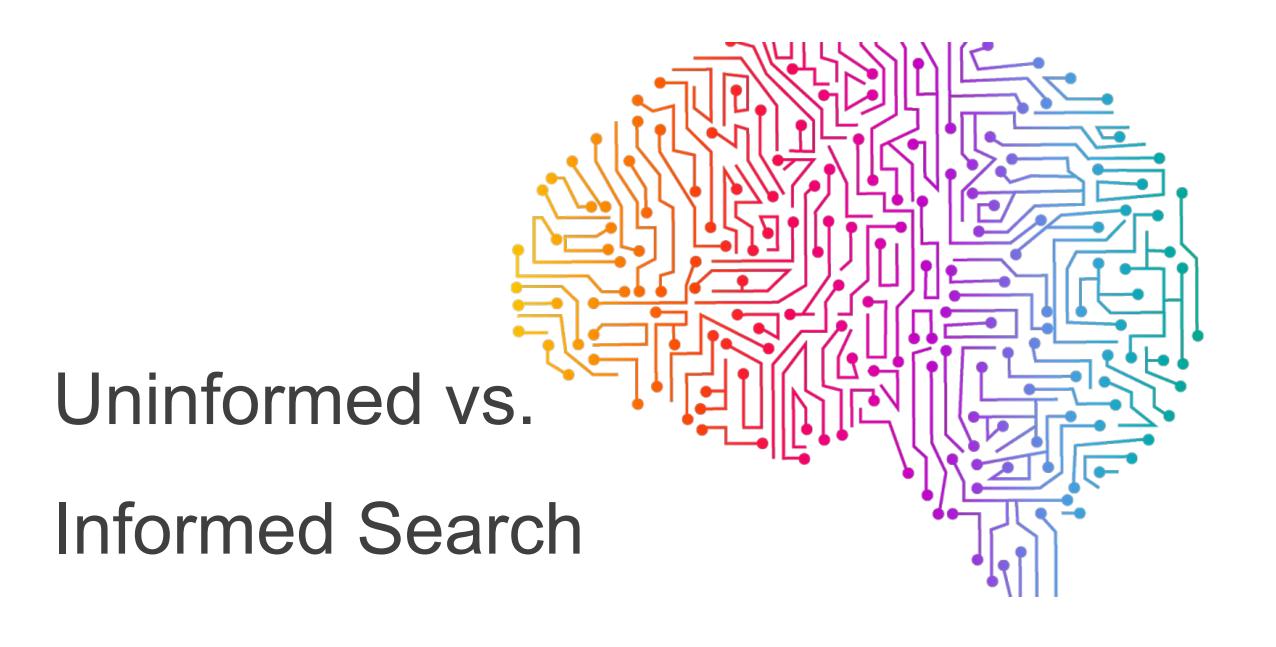
Week #2-2

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Learning Objective of this Topic

- Uninformed vs. Informed Search
- Uninformed Searching Algorithms
 - ☐ Breadth First Search (BFS)
 - Depth First Search (DFS)
 - Uninformed Cost Search (UCS)
 - Difference between BFS and DFS
 - ☐ When to use BFS? Real-Life Applications
 - When to use DFS? Real-Life Applications



Uninformed vs. Informed Search

Uninformed

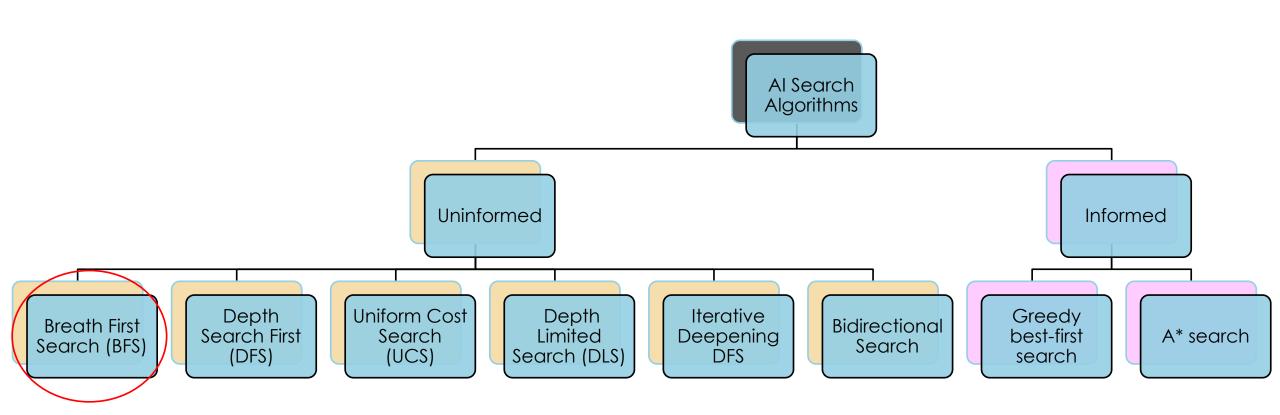
- Also known as Blind search.
- Do not require information to Derform a search.
- May be time-consuming.
- Comparatively high in cost.
- The Al does not get any suggestions regarding the solution and where to find it.

Informed

- Also known as Heuristic search.
- Requires information to performsearch Quick solution to the problem.
- Cost is low.
- The Al gets suggestions regarding how and where to find a solution to any problem.



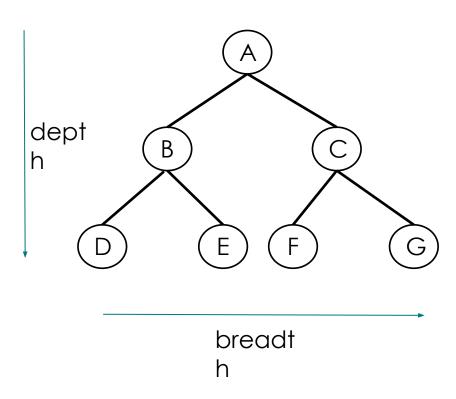
UNINFORMED SEARCH ALGORITHM



Breadth First

Search (BFS)

- Expand Shallowest Node First (Level First Search)
- Queue: nodes in the queue to be explored
- explored: Nodes that are already explored
- Queue is a first-in-first-out (FIFO) queue, i.e., new successors go at the end of the queue.



- Expand Shallowest Node First (Level First Search)
- Queue: nodes in the queue to be explored
- explored: Nodes that are already explored
- Queue is a first-in-first-out (FIFO) queue, i.e., new successors go at the end of the queue.

```
function BREADTH-FIRST-SEARCH(problem) returns a solution, or failure

Queue ← a FIFO queue with source node as the only element

explored ← an empty set

if source node == goal then return SOLUTION

loop do

if EMPTY?( Queue ) then return failure

node ← POP( Queue ) /* chooses the shallowest node in Queue */

add node to explored

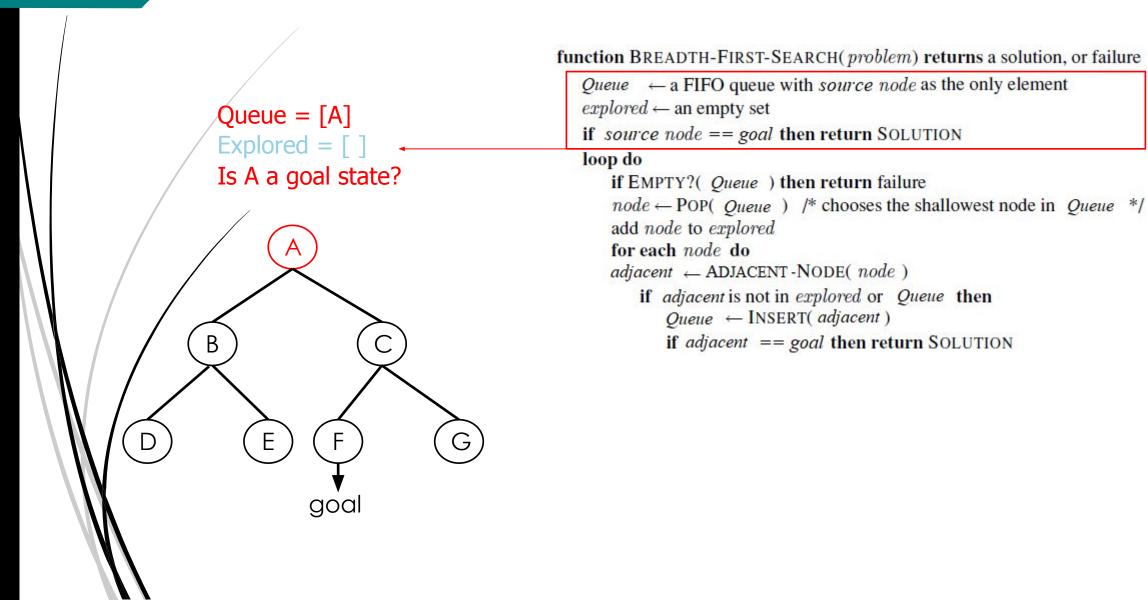
for each node do

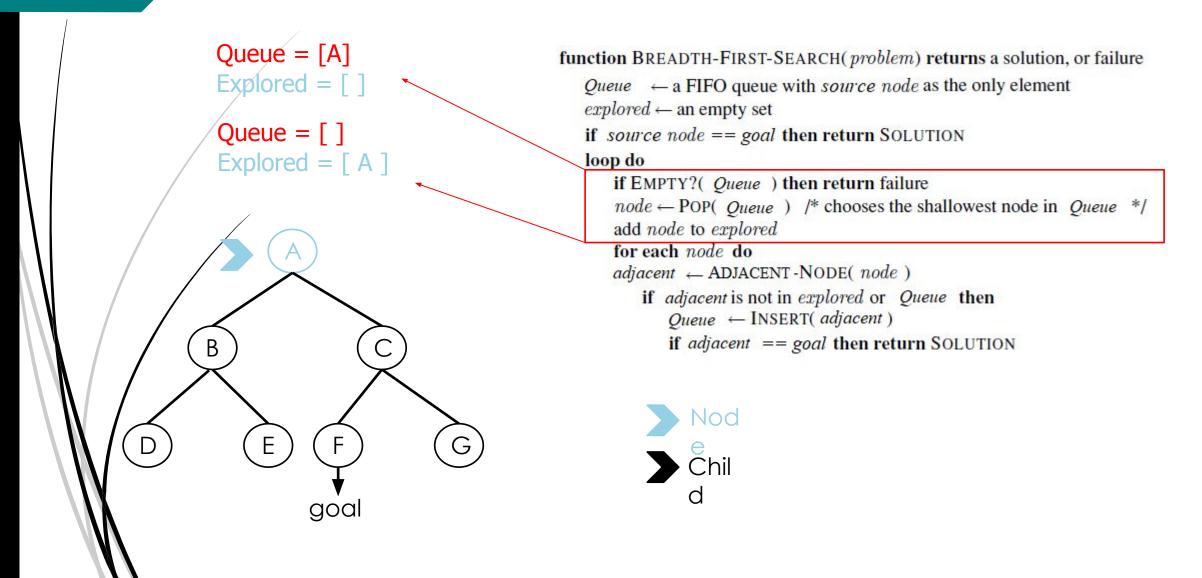
adjacent ← ADJACENT-NODE( node )

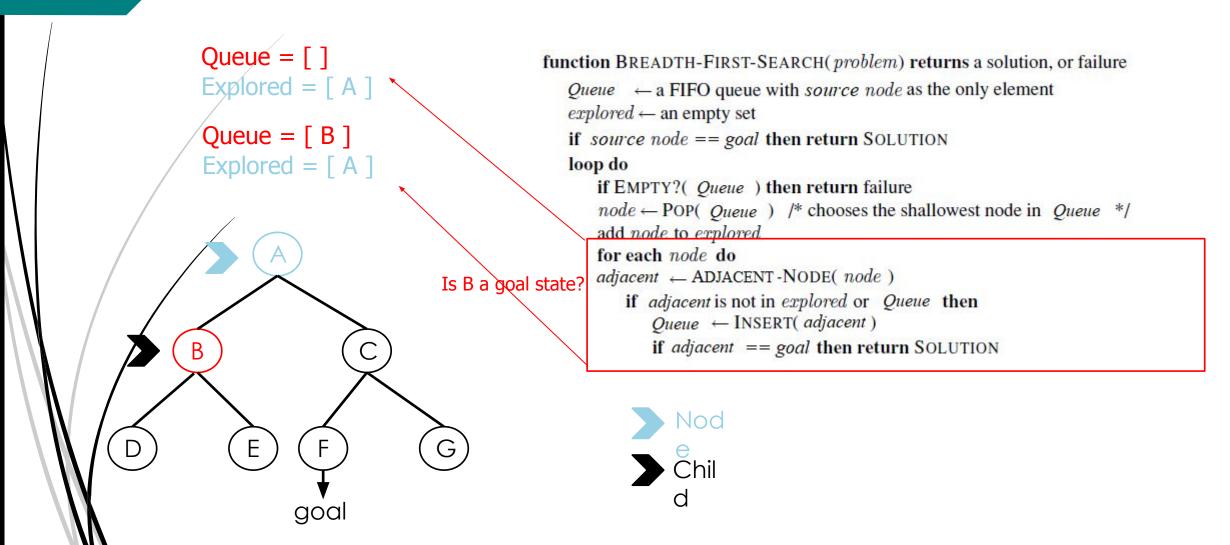
if adjacent is not in explored or Queue then

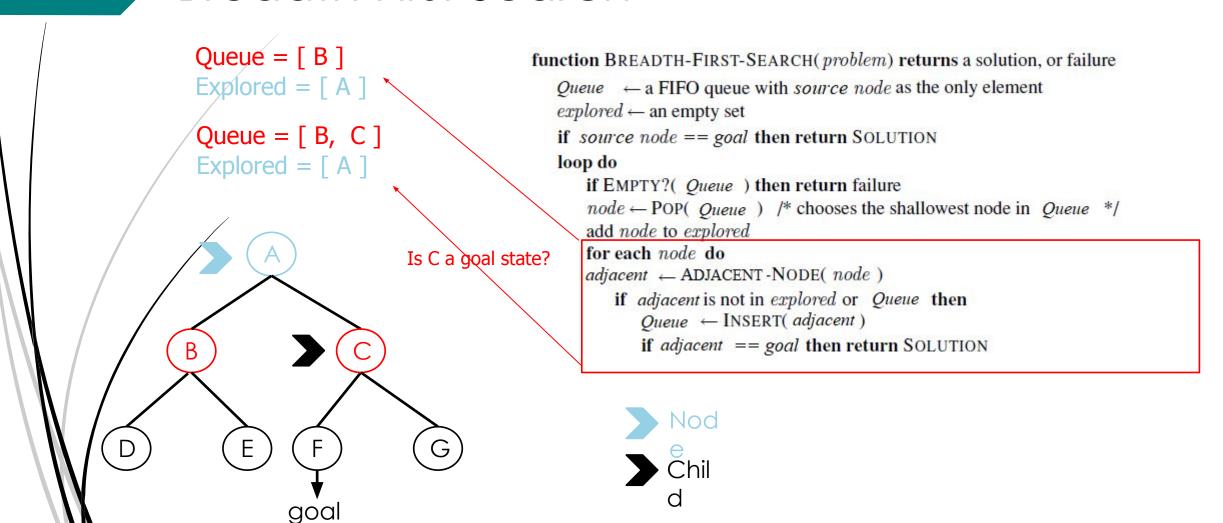
Queue ← INSERT( adjacent )

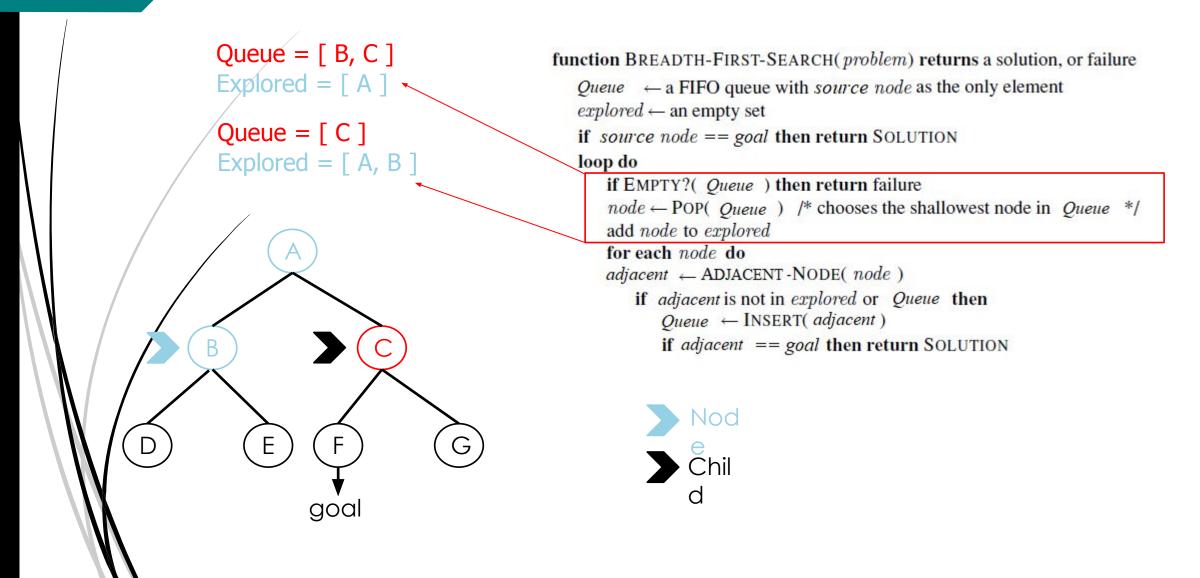
if adjacent == goal then return SOLUTION
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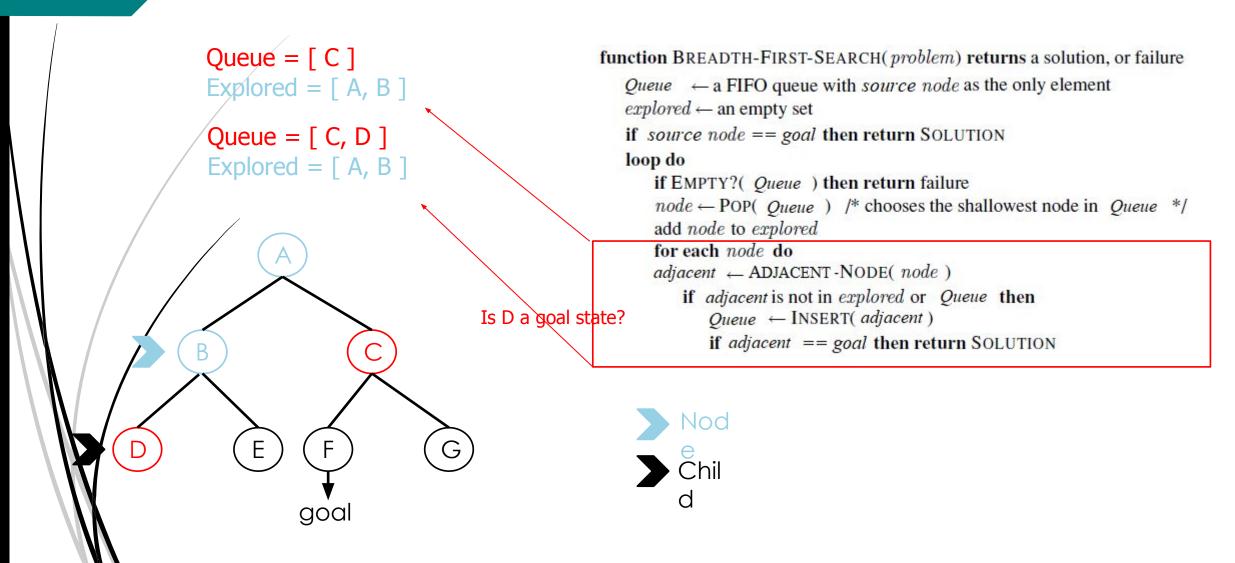


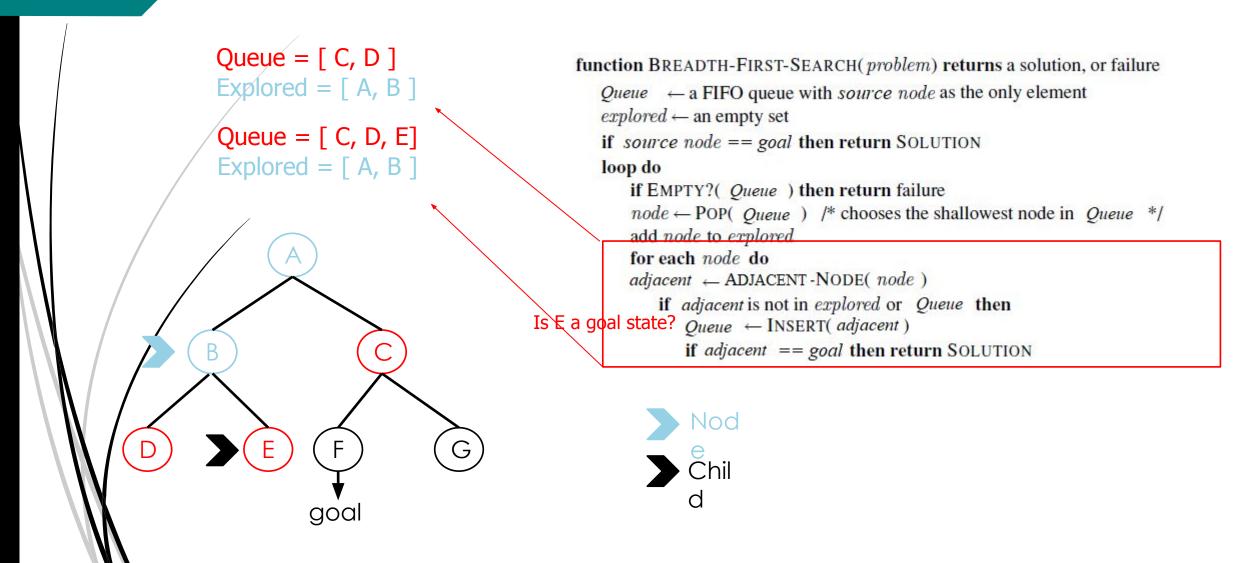


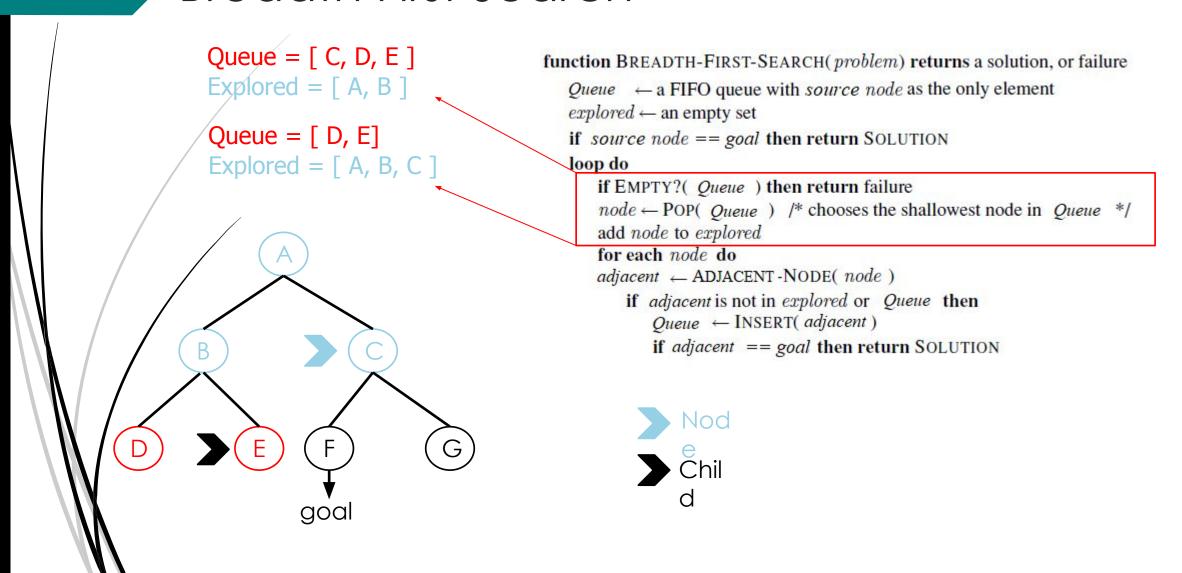


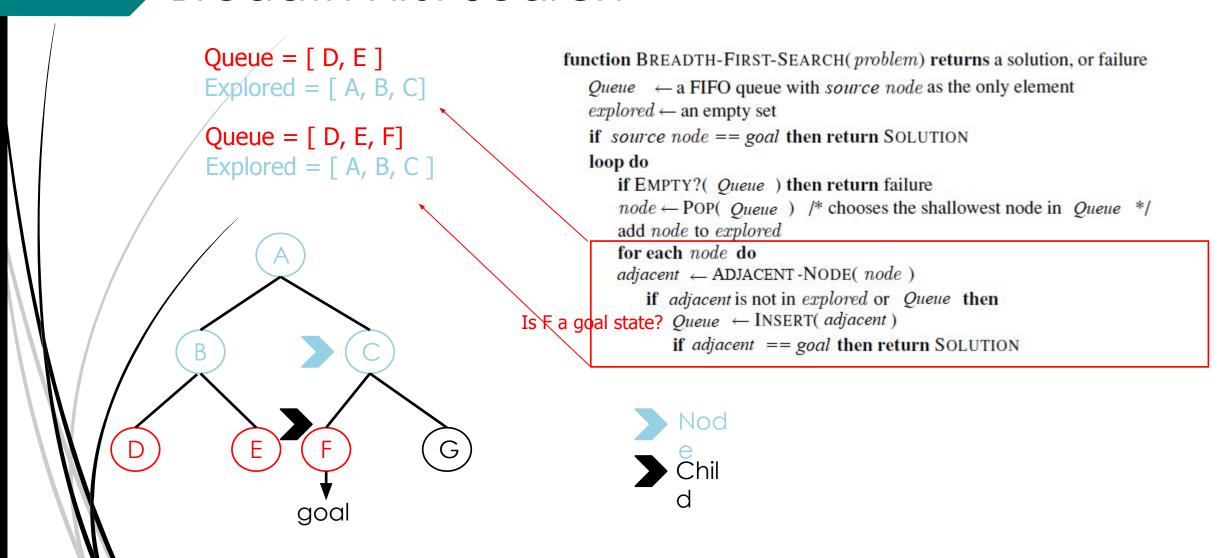




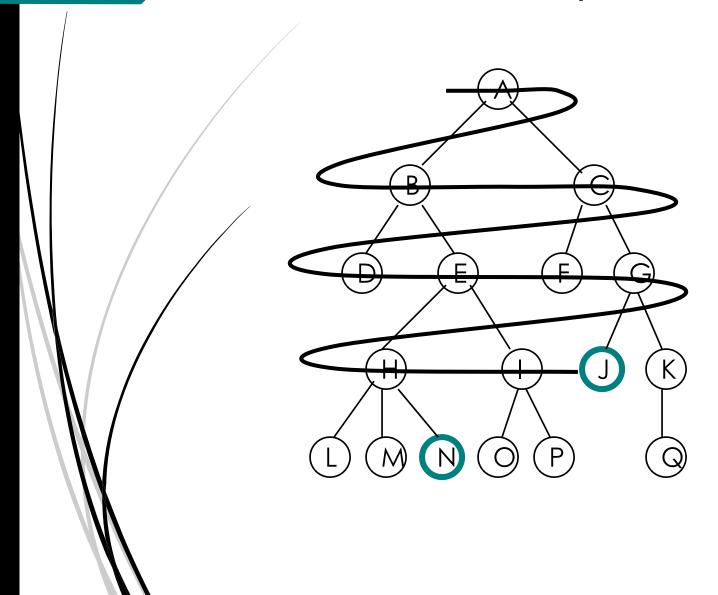








BFS: Summary



BFS: Completeness

- Completeness: Yes (guaranteed to find a solution if there exists one)
- if the shallowest goal node is at some finite depth d, breadth-first search will eventually find it after generating all shallower nodes (provided the branching factor b is finite).

the branching factor is the number of children at each node, the outdegree.

BFS: Optimal?

- Not necessarily optimal
- Only optimal if every action has same cost.

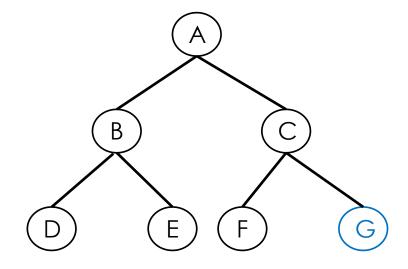
BFS: Time Complexity?

- Worst complexity is when G is a goal state.
- In this case, total number of nodes generated is

$$b+b^2+b^3+...+b^d=O(b^d)$$

(The dth layer contains nodes much larger than all the nodes in previous layers combined!)

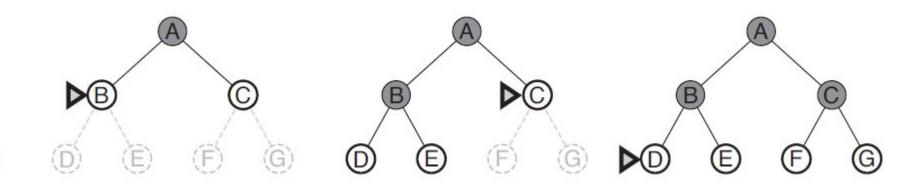
So the time complexity is O(bd)

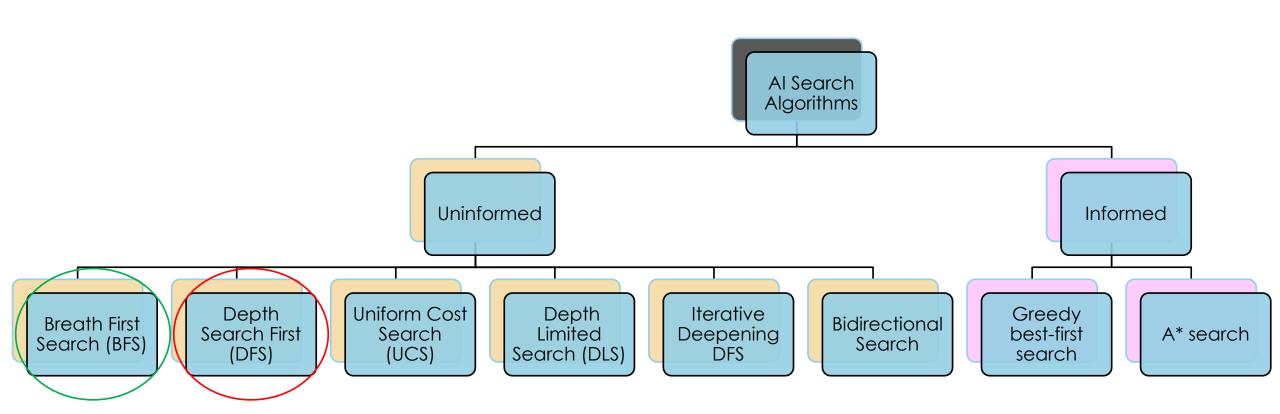


A binary tree, b=2, d=2

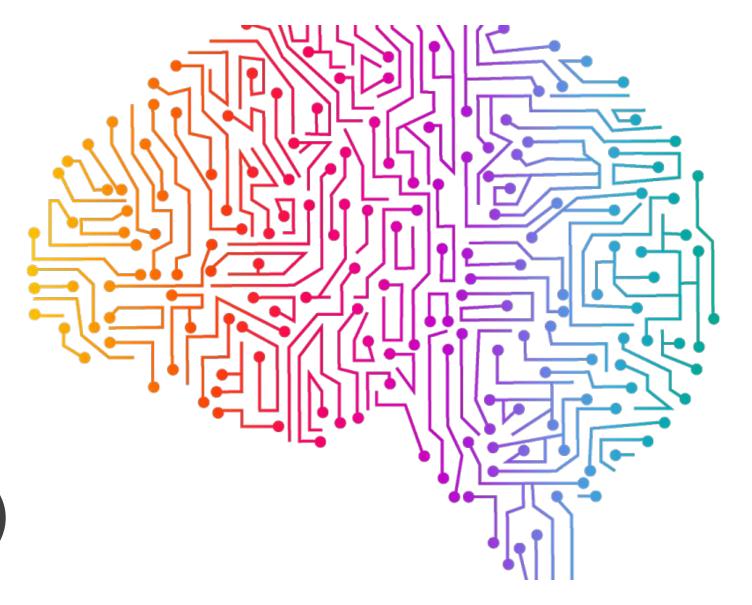
BFS: Space Complexity?

- There will be $O(b^{d-1})$ nodes in the explored set and $O(b^d)$ nodes in the queue.
- \square So the space complexity is $O(b^d)$, i.e., it is dominated by the size of the queue.
- Exponential time complexity can be accepted but exponential space complexity is BAD!



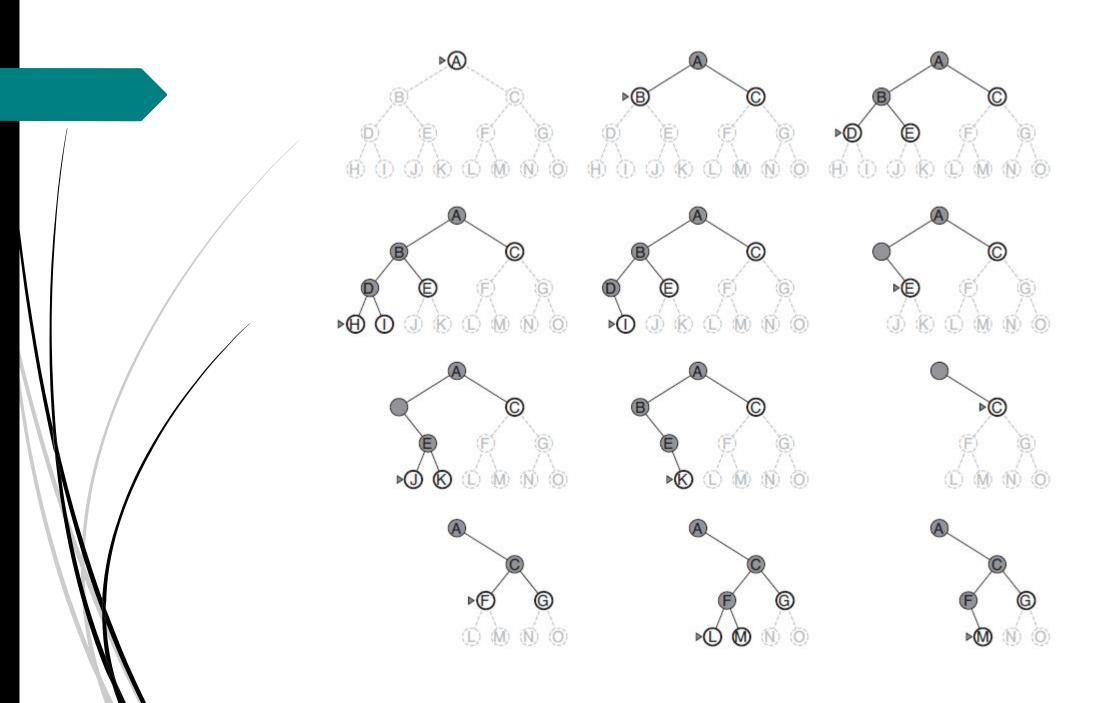


Depth First
Search (DFS)



Depth First Search

- ☐ It works vertically.
- ☐ It doesn't generate a solution always.
- It involves exhaustive searches of all the nodes by going ahead, if possible, else by backtracking.
- Here, the word backtrack means that when you are moving forward and there are no more nodes along the current path, you move backward on the same path to find nodes to traverse.



M is the goal node

Depth First Search

Expand Depth Node First

(Depth First Search)

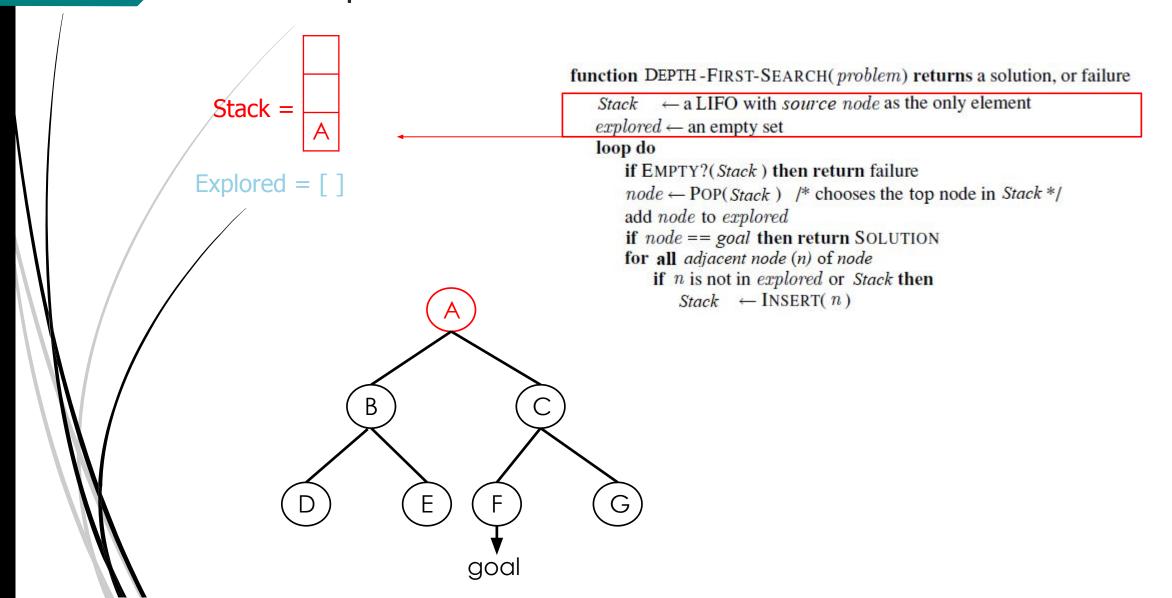
- Stack: nodes in the stack to be explored
- explored: Nodes that are already explored
- Stack is a Last-in-first-out (LIFO) queue.

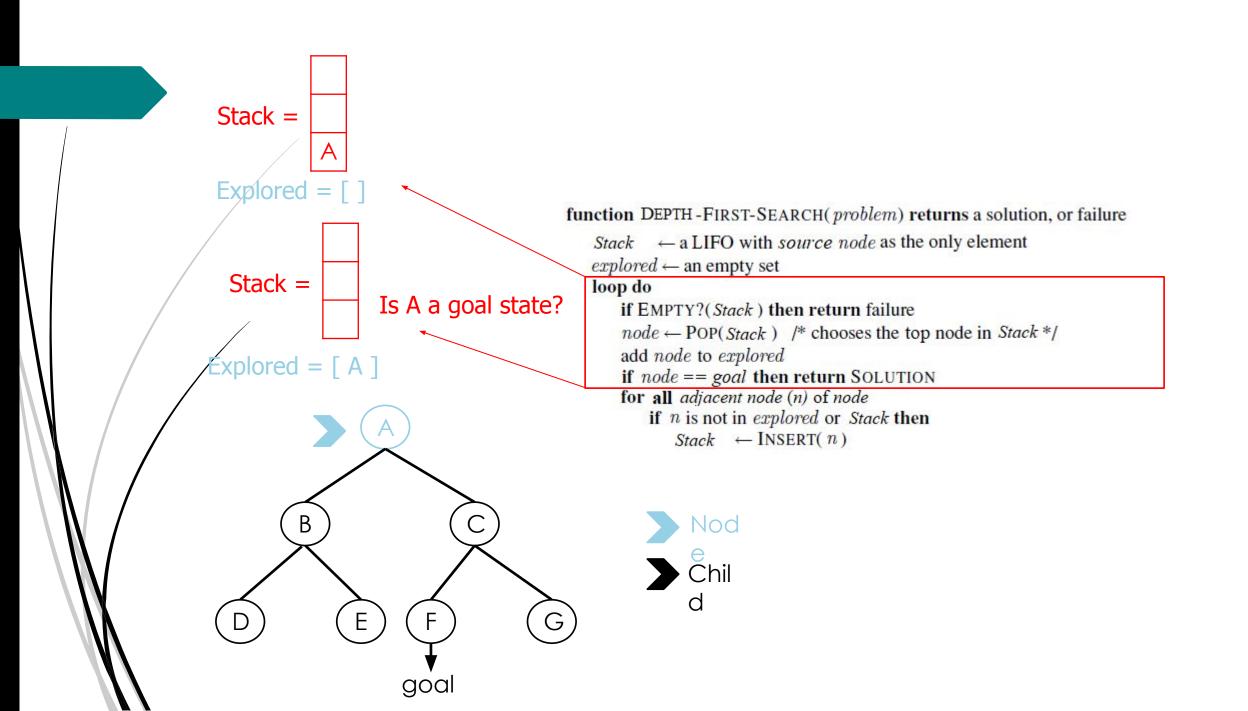
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function DEPTH-FIRST-SEARCH(problem) returns a solution, or failure

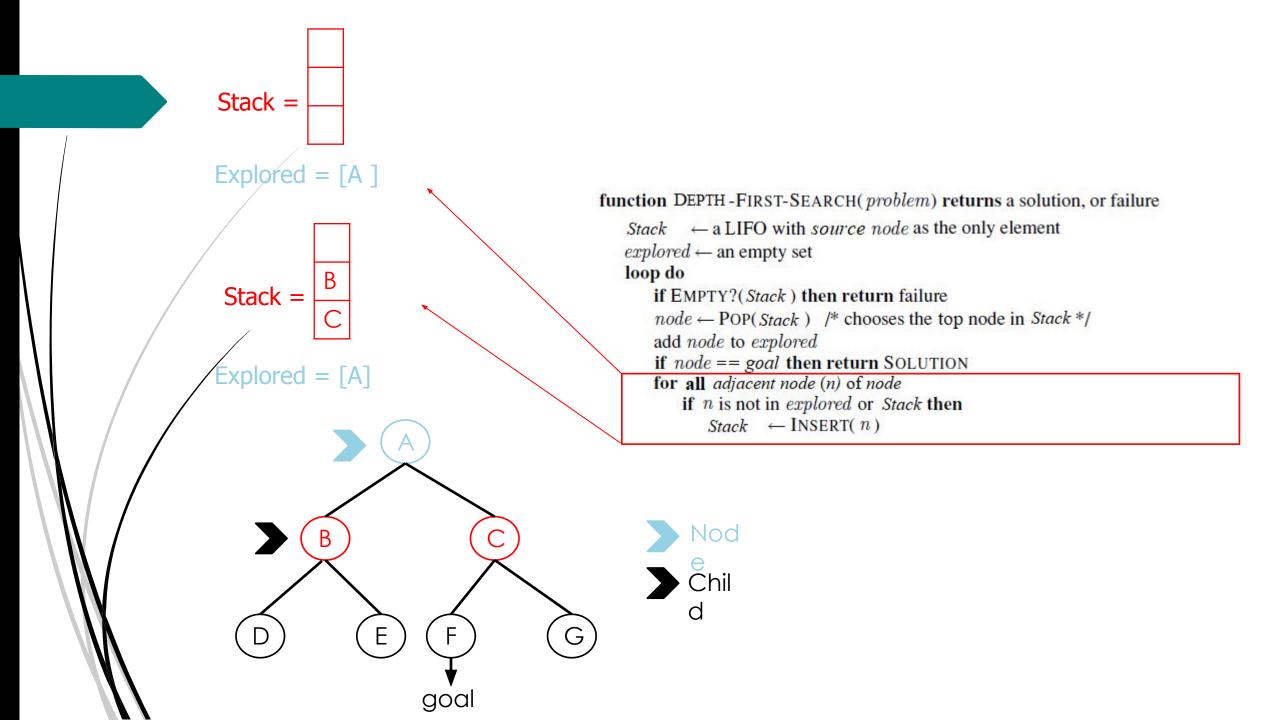
Stack ← a LIFO with source node as the only element
explored ← an empty set
loop do

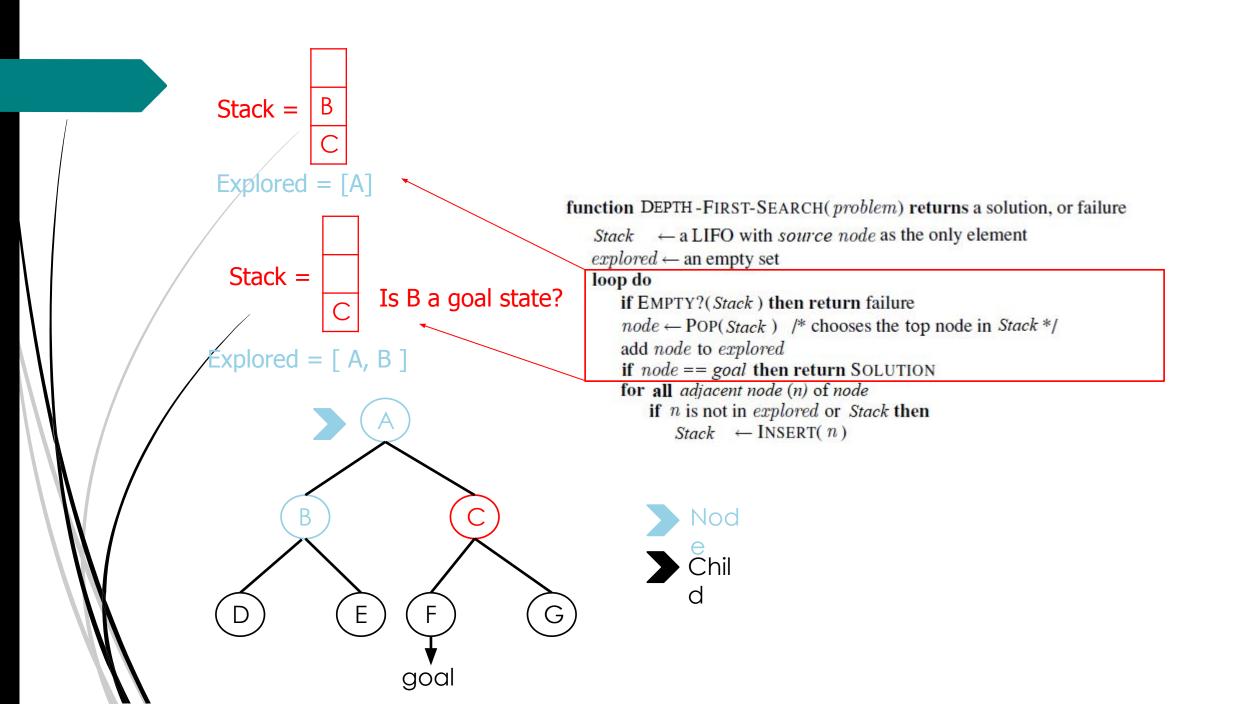
if EMPTY?(Stack) then return failure
node ← POP(Stack) /* chooses the top node in Stack*/
add node to explored
if node == goal then return SOLUTION
for all adjacent node (n) of node
if n is not in explored or Stack then
Stack ← INSERT(n)
```

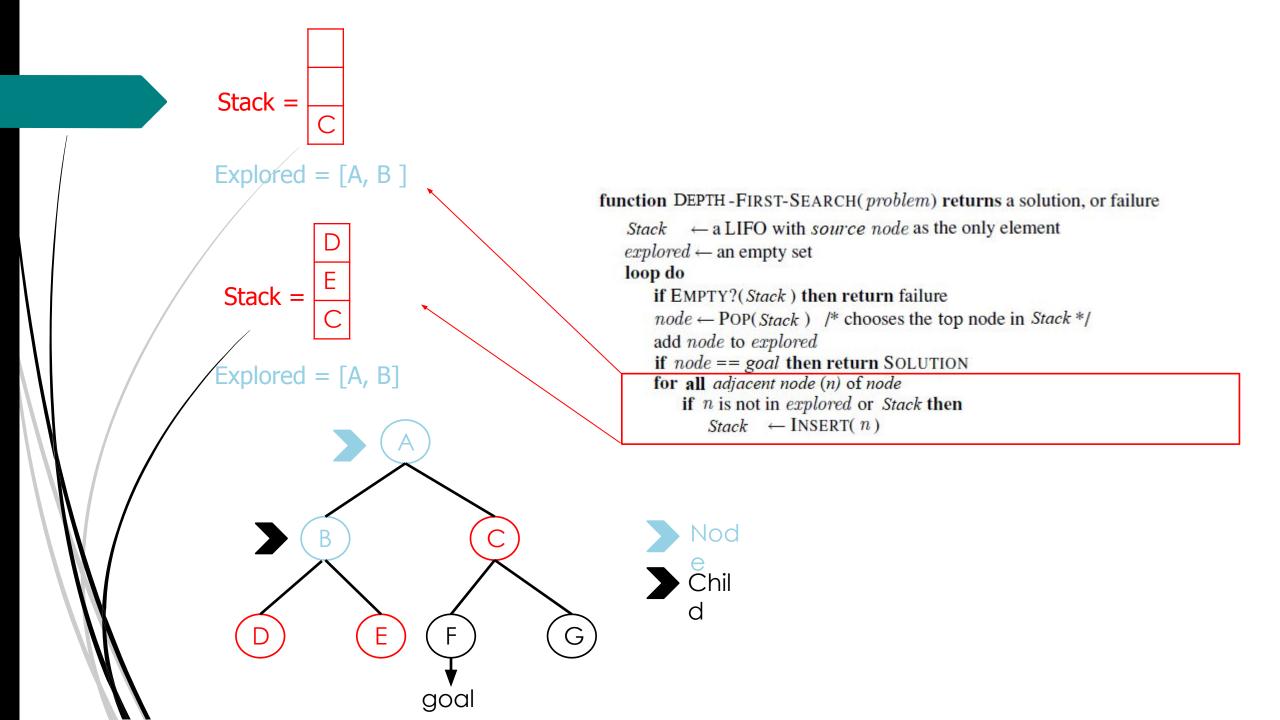
Depth First Search

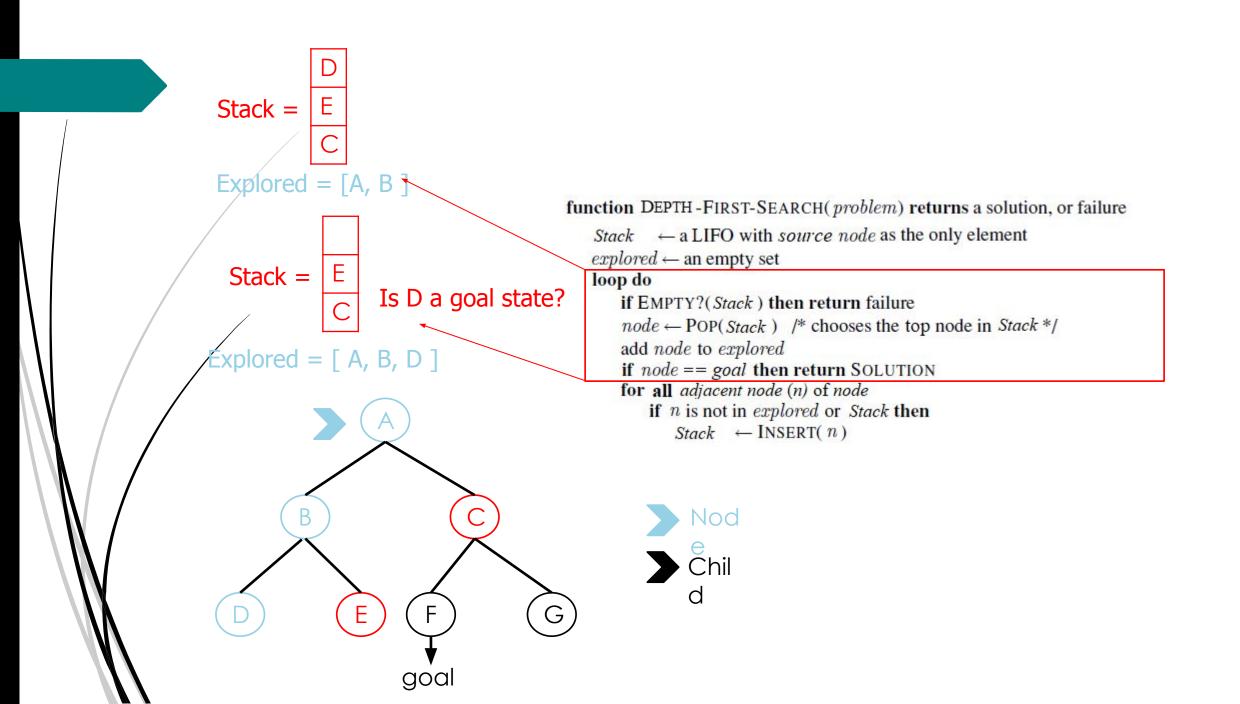


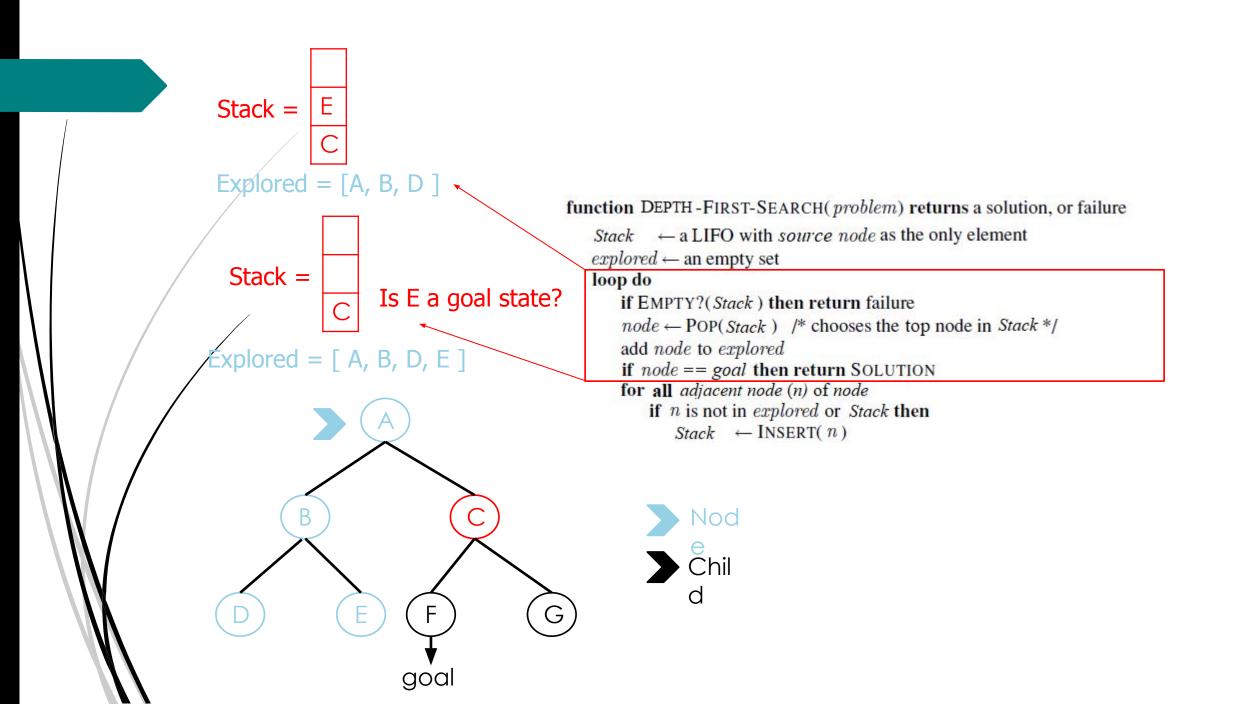


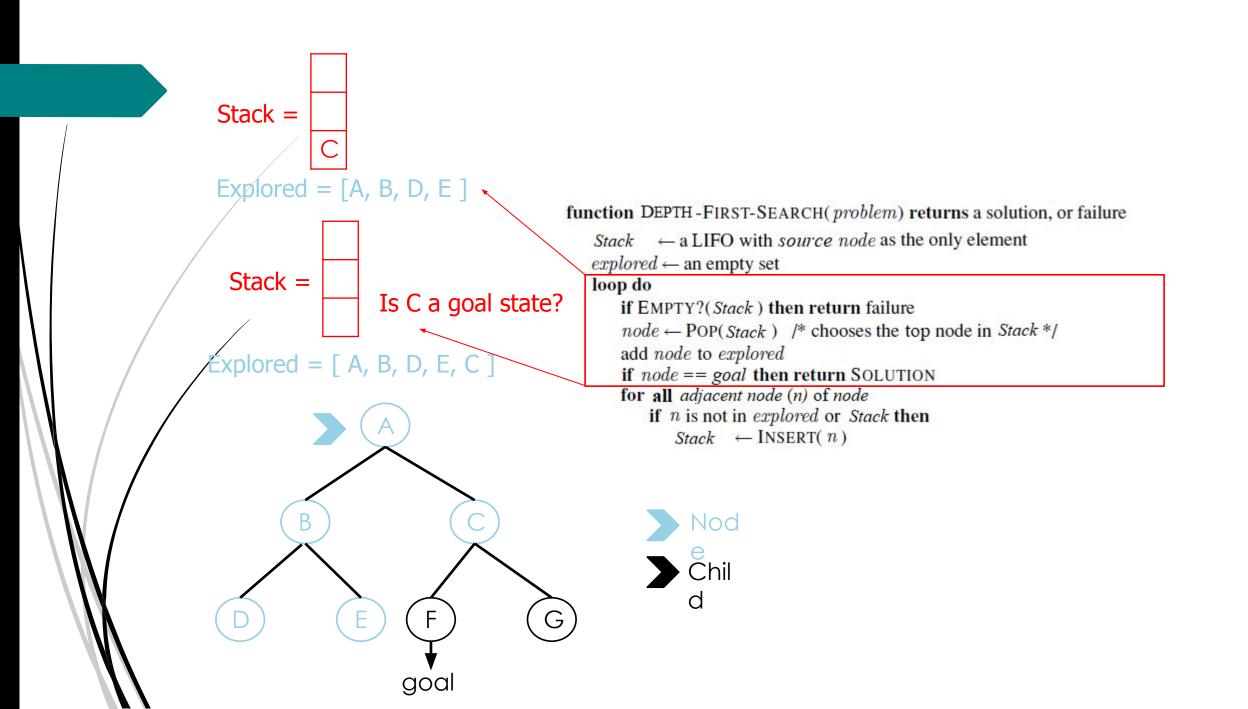


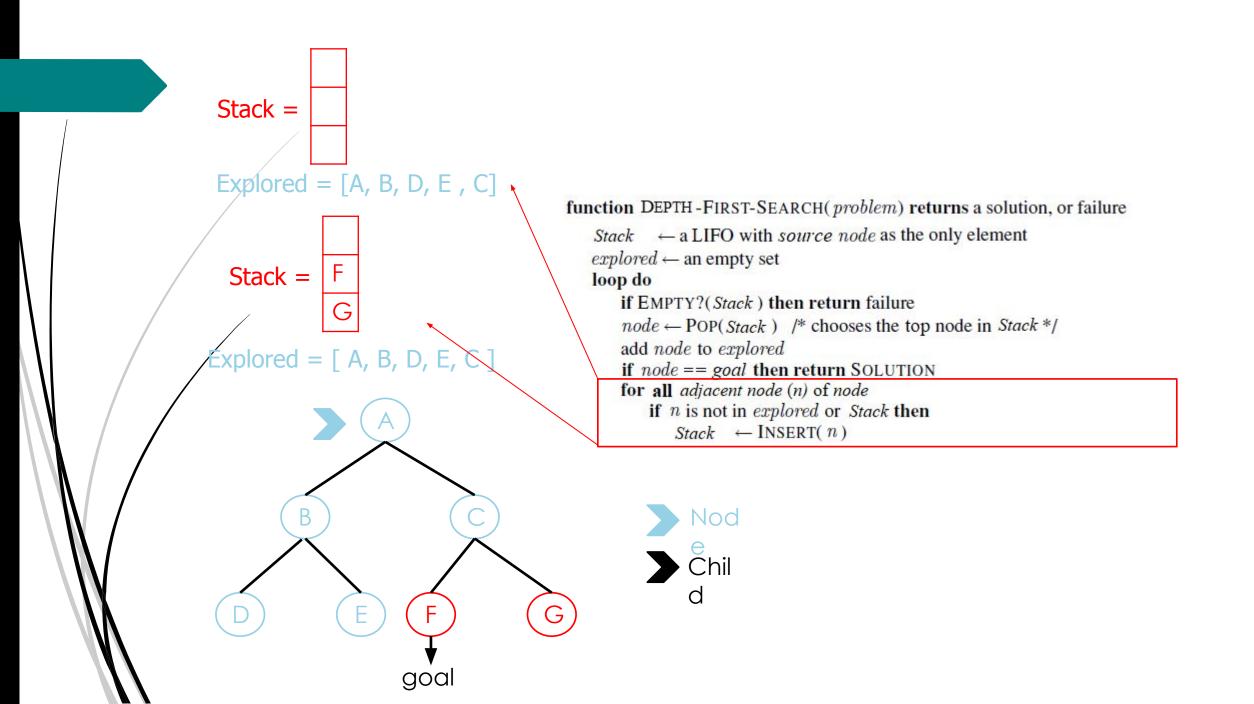


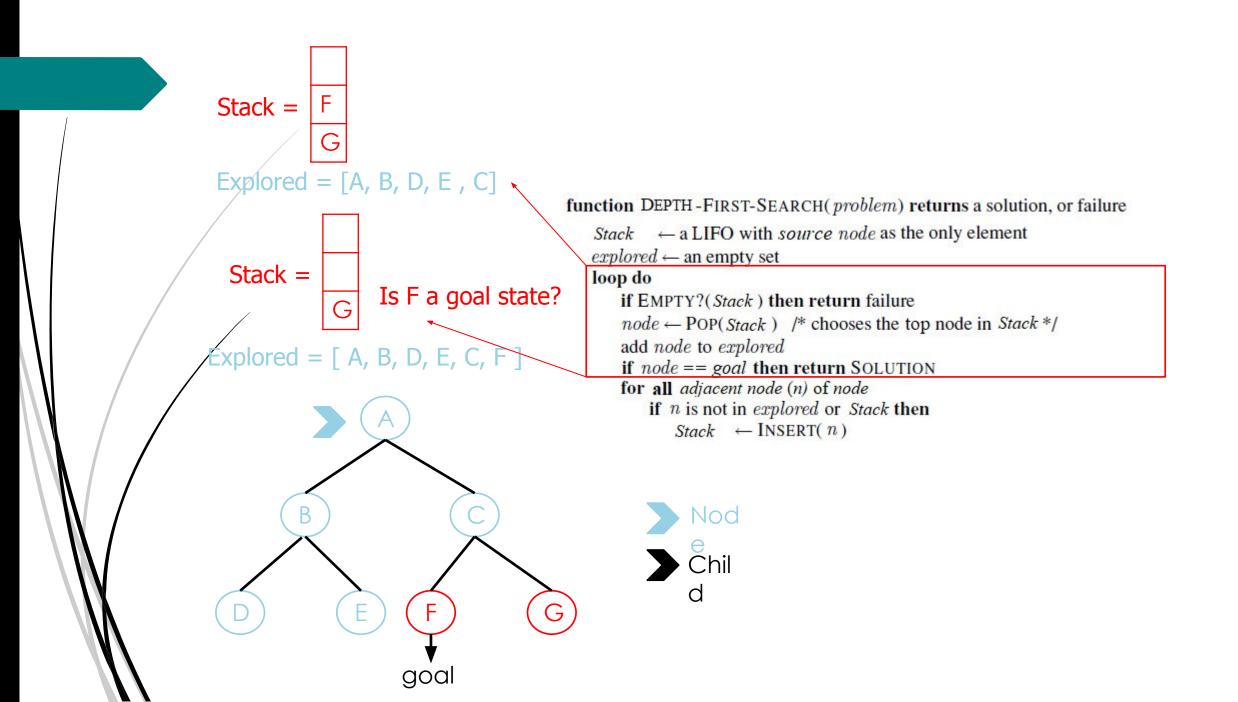






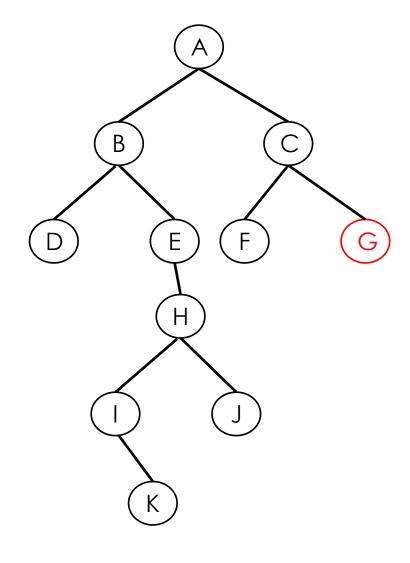






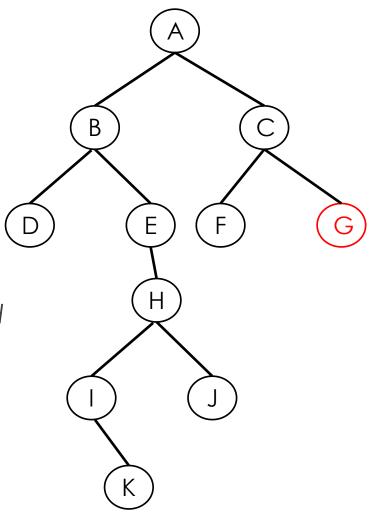
Completeness?

- Goal, G is at depth, d=2
- \square Max depth of tree, m = 5
- DFS is complete ONLY if m is finite



Time Complexity

- Number of nodes visited? O(b^m)
- For BFS, remember it was O(bd)
- \square DFS is terrible if m is much larger than d
 - but if solutions are dense (goal is at K instead of G), DFS is faster than BFS.

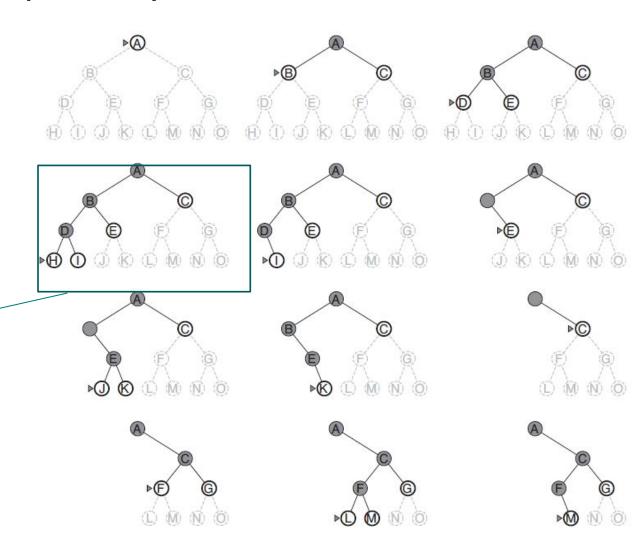


Space Complexity?

(bm), i.e., linear space!

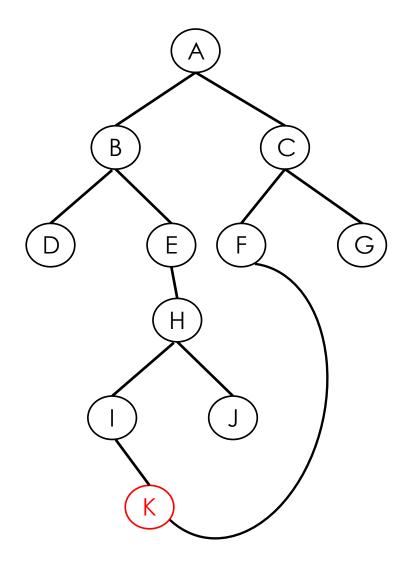
(we only need to remember a single path + expanded unexplored nodes)

O(bm)? E.g. m=3, b=2 (we need to save max of 3*2=6 nodes in memory)



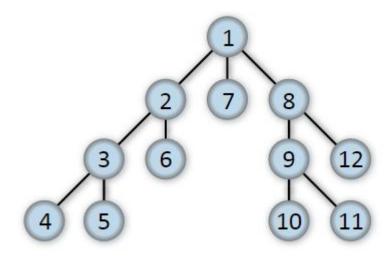
Optimal?

- Remember BFS was optimal if each step has same cost.
- Eyen if this is the case, DFS is not guaranteed to find the optimum path to the goal.
- For the graph on the right, BFS will return optimal path but DFS won't (assuming each step has same cost)



Class Activity

Find the Goal Node 9 with the help of DFS using Stack.



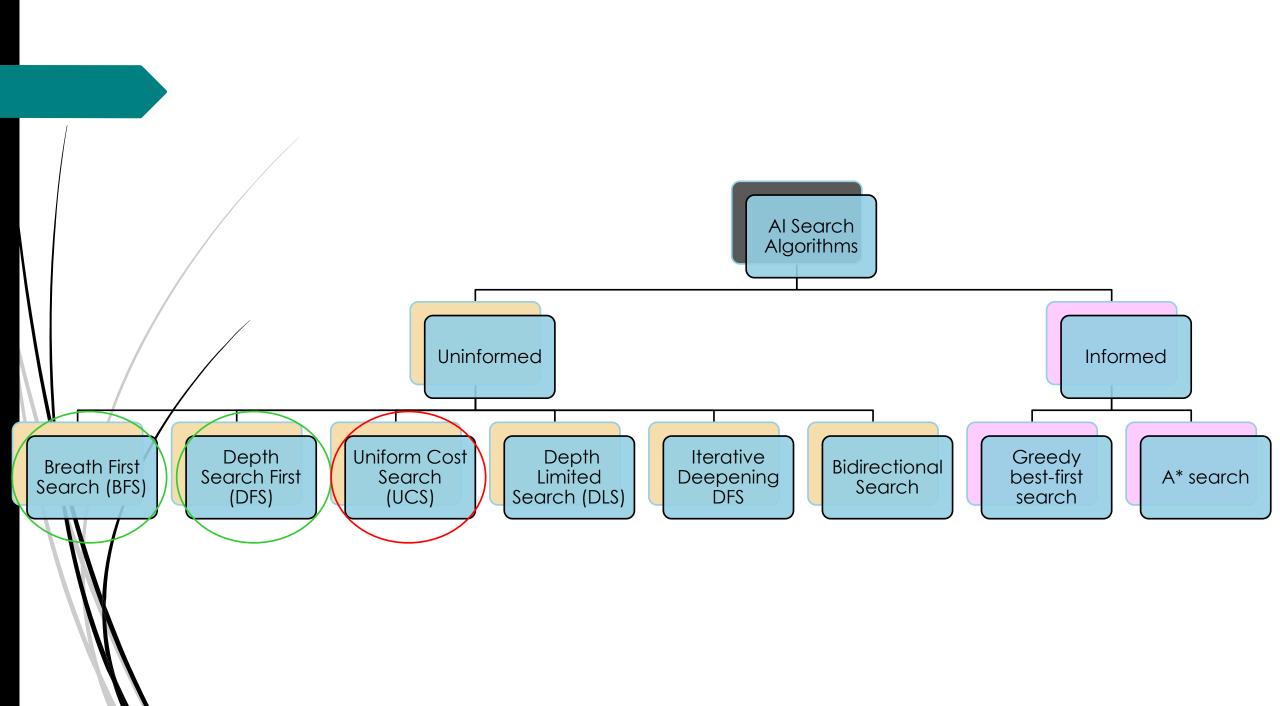
Difference between BFS and DFS

BFS

- Visit nodes level by level.
- Uses Queue data structure.
- May give an optimal solution.
- ☐ Slower, need more memory.

DFS

- ☐ Visit nodes graph depth-wise.
- Uses Stack data structure.
- Do not guarantee an optimal solution.
- ☐ Faster, need less memory.



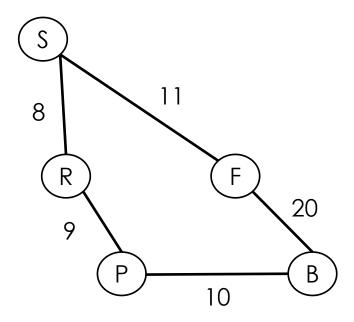


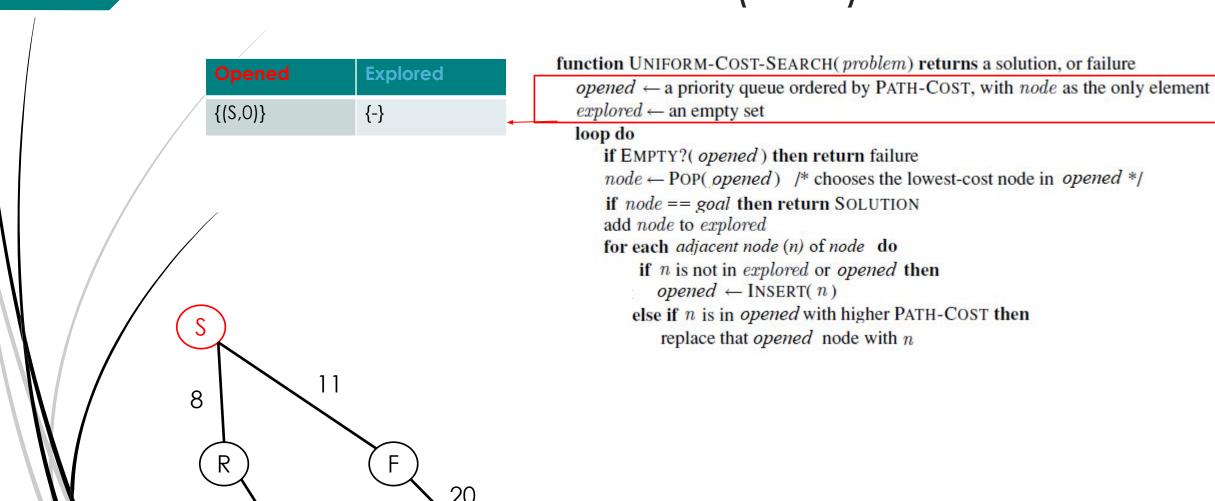
Uniform Cost Search (UCS) --- How it is different from BFS and DFS

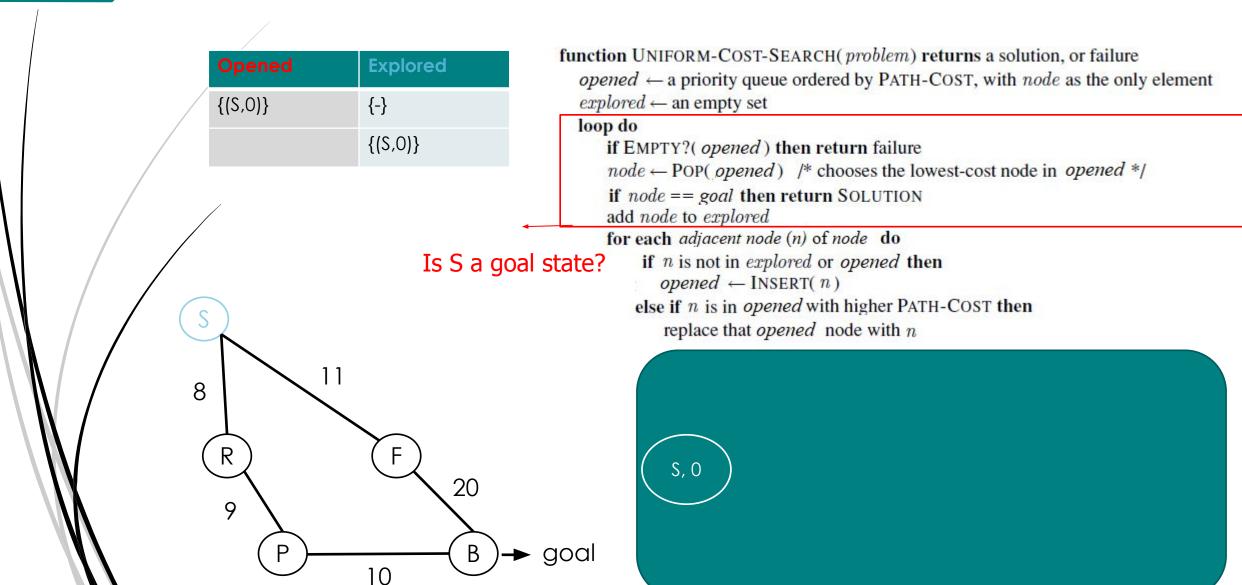
- In contrast to BFS and DFS algorithms that don't take into consideration either the cost between two nodes, the UCS algorithm uses the path's cost from the initial node to the current node as the extension criterion.
- 1 Starting from the initial state (starting node), the UCS algorithm, in each step chooses the node that is closer to the initial node.
- When the algorithm finds the solution, returns the path from the initial state to the final state.
- The UCS algorithm is characterized as complete, as it always returns a solution if exists.
- Moreover, the UCS algorithm guarantees the optimum solution.

- Uniform cost search modifies BFS such that it works with any cost function.
- Instead of expanding the shallowest node, uniform-cost search expands the node n with the lowest path cost.
- This is done by storing the opened list as a priority queue ordered by cost.

- Expand Node with lowest path cost
- Opened: a priority queue ordered by path cost
- explored: Nodes that are already explored







Opened	Explored
{(S,O)}	{-}
{(R, 8), (F, 11)}	{(S,O)}

function UNIFORM-COST-SEARCH(problem) returns a solution, or failure
 opened ← a priority queue ordered by PATH-COST, with node as the only element
 explored ← an empty set

loop do

if EMPTY?(opened) then return failure

 $node \leftarrow POP(opened)$ /* chooses the lowest-cost node in opened */

if node == goal then return SOLUTION

add node to explored

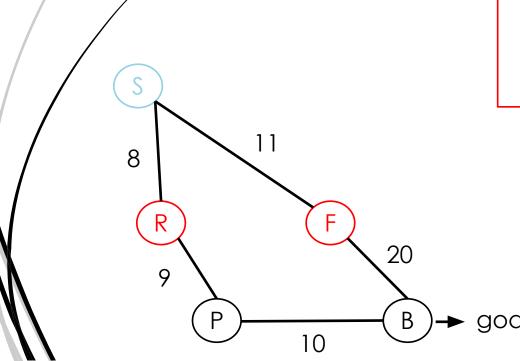
for each adjacent node (n) of node do

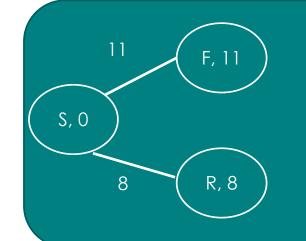
if n is not in explored or opened then

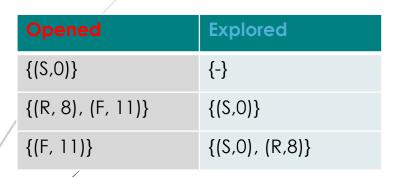
 $opened \leftarrow INSERT(n)$

else if n is in *opened* with higher PATH-COST then

replace that opened node with n







function UNIFORM-COST-SEARCH(problem) returns a solution, or failure $opened \leftarrow$ a priority queue ordered by PATH-COST, with node as the only element $explored \leftarrow$ an empty set

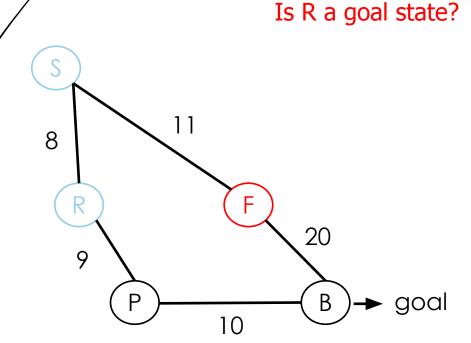
loop do

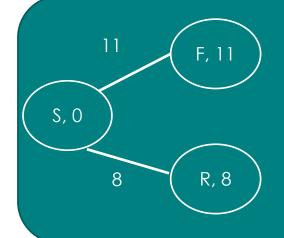
if EMPTY?(opened) then return failure
node ← POP(opened) /* chooses the lowest-cost node in opened */
if node == goal then return SOLUTION
add node to explored

for each adjacent node (n) of node do

if n is not in explored or opened then opened \leftarrow INSERT(n)

else if n is in *opened* with higher PATH-COST then replace that *opened* node with n





Opened	Explored
{(0,2)}	{-}
{(R, 8), (F, 11)}	{(S,O)}
{(F, 11)}	{(S,0), (R,8)}
{(F, 11), (P, 17)}	{(S,0), (R,8)}

function UNIFORM-COST-SEARCH(problem) returns a solution, or failure
 opened ← a priority queue ordered by PATH-COST, with node as the only element
 explored ← an empty set
loop do

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if node == goal then return SOLUTION

add node to explored

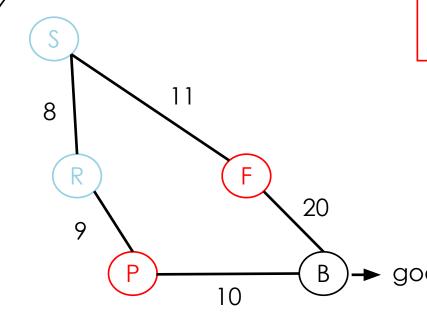
for each adjacent node (n) of node do

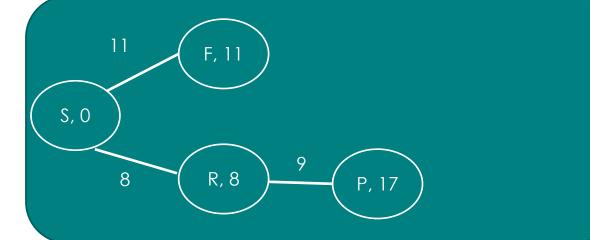
if n is not in explored or opened then

 $opened \leftarrow INSERT(n)$

else if n is in *opened* with higher PATH-COST then

replace that opened node with n





Opened	Explored
{(S,O)}	{-}
{(R, 8), (F, 11)}	{(S,O)}
{(F, 11)}	{(S,0), (R,8)}
{(F, 11), (P, 17)}	{(S,0), (R,8)}
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function UNIFORM-COST-SEARCH(problem) returns a solution, or failure $opened \leftarrow$ a priority queue ordered by PATH-COST, with node as the only element $explored \leftarrow$ an empty set

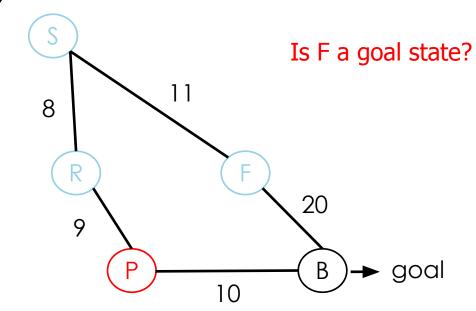
loop do

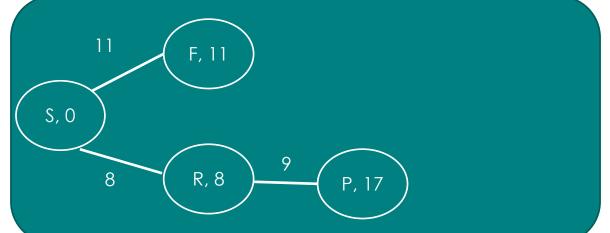
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node ← POP(opened) /* chooses the lowest-cost node in opened */
if node == goal then return SOLUTION

add node to explored for each adjacent node (n) of node do

if n is not in explored or opened then opened \leftarrow INSERT(n)

else if n is in *opened* with higher PATH-COST then replace that *opened* node with n



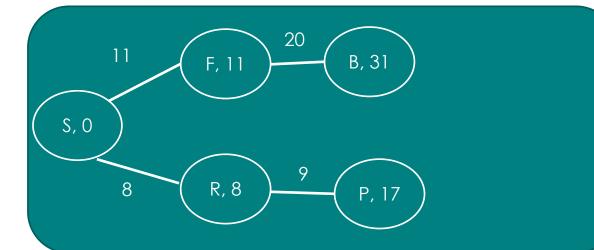


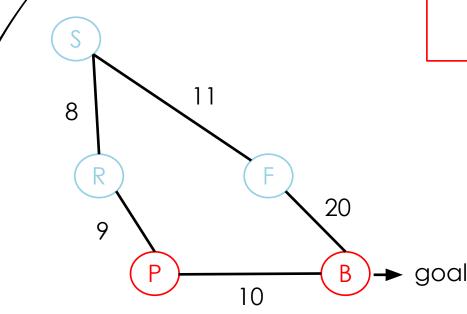
Opened	Explored
{(S,O)}	{-}
{(R, 8), (F, 11)}	{(S,O)}
{(F, 11)}	{(S,0), (R,8)}
{(F, 11), (P, 17)}	{(S,0), (R,8)}
{(P, 17)}	{(S,0), (R,8), (F, 11)}
{(P, 17), (B, 31)}	{(S,0), (R,8), (F, 11)}

function UNIFORM-COST-SEARCH(problem) returns a solution, or failure
 opened ← a priority queue ordered by PATH-COST, with node as the only element
 explored ← an empty set
loop do
 if EMPTY?(opened) then return failure
 node ← POP(opened) /* chooses the lowest cost node in opened */

node ← POP(opened) /* chooses the lowest-cost node in opened */
if node == goal then return SOLUTION
add node to explored

for each adjacent node (n) of node do
 if n is not in explored or opened then
 opened ← INSERT(n)
 else if n is in opened with higher PATH-COST then
 replace that opened node with n





Opened Explored

{(0,0)}	{-}
{(R, 8), (F, 11)}	{(S,O)}
{(F, 11)}	{(S,0), (R,8)}
{(F, 11), (P, 17)}	{(S,0), (R,8)}
{(P, 17)}	{(S,0), (R,8), (F, 11)}
{(P, 17), (B, 31)}	{(S,0), (R,8), (F, 11)}
{(B, 31)}	{(S,0), (R,8), (F, 11),(P, 17)}

function UNIFORM-COST-SEARCH(problem) **returns** a solution, or failure $opened \leftarrow$ a priority queue ordered by PATH-COST, with node as the only element $explored \leftarrow$ an empty set

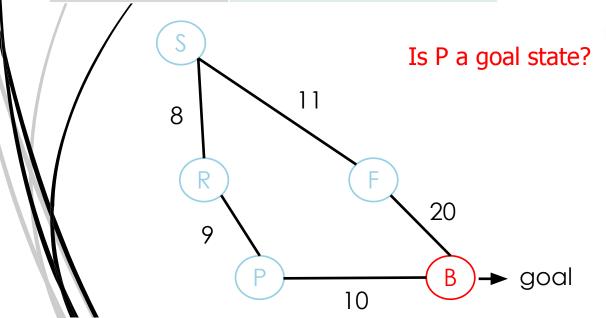
loop do

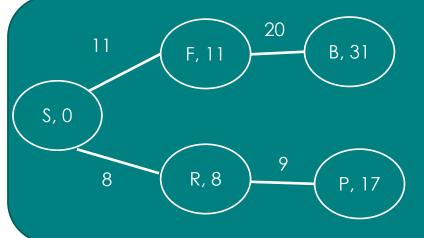
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else if n is in *opened* with higher PATH-COST then replace that *opened* node with n

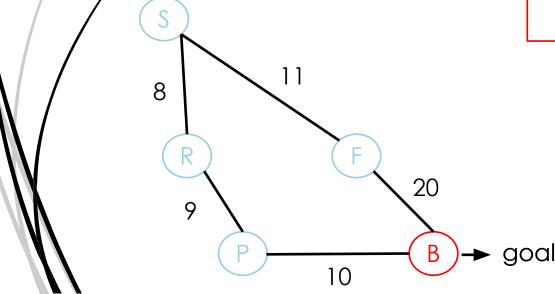


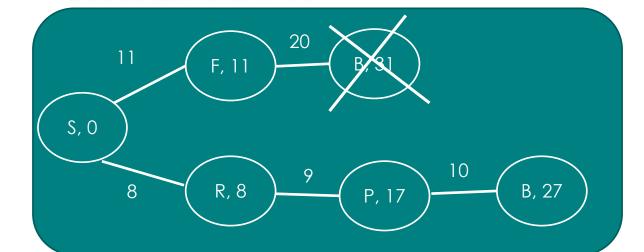


Opened	Explored
{(0,0)}	{-}
{(R, 8), (F, 11)}	{(0,0)}
{(F, 11)}	{(S,0), (R,8)}
{(F, 11), (P, 17)}	{(S,0), (R,8)}
{(P, 17)}	{(S,0), (R,8), (F, 11)}
{(P, 17), (B, 31)}	{(S,0), (R,8), (F, 11)}
{(B, 31)}	{(S,0), (R,8), (F, 11),(P, 17)}
{(B, 27)}	

function UNIFORM-COST-SEARCH(problem) returns a solution, or failure
 opened ← a priority queue ordered by PATH-COST, with node as the only element
 explored ← an empty set
loop do
 if EMPTY?(opened) then return failure
 node ← POP(opened) /* chooses the lowest-cost node in opened */
 if node == goal then return SOLUTION
 add node to explored

for each adjacent node (n) of node do
 if n is not in explored or opened then
 opened ← INSERT(n)
 else if n is in opened with higher PATH-COST then
 replace that opened node with n





Opened	Explored
{(S,O)}	{-}
{(R, 8), (F, 11)}	{(S,O)}
{(F, 11)}	{(S,0), (R,8)}
{(F, 11), (P, 17)}	{(S,O), (R,8)}
{(P, 17)}	{(S,O), (R,8), (F, 11)}
{(P, 17), (B, 31)}	{(S,O), (R,8), (F, 11)}
{(B, 31)}	{(S,0), (R,8), (F, 11),(P, 17)}
{(B, 27)}	

function UNIFORM-COST-SEARCH(problem) **returns** a solution, or failure $opened \leftarrow$ a priority queue ordered by PATH-COST, with node as the only element $explored \leftarrow$ an empty set

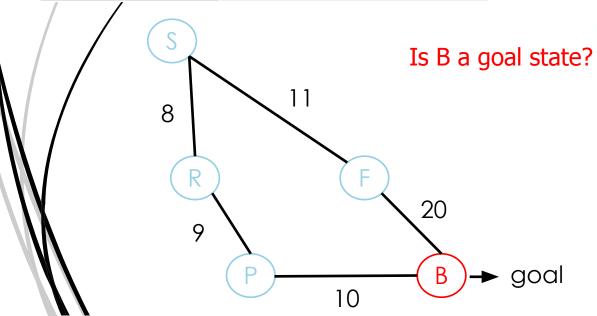
loop do

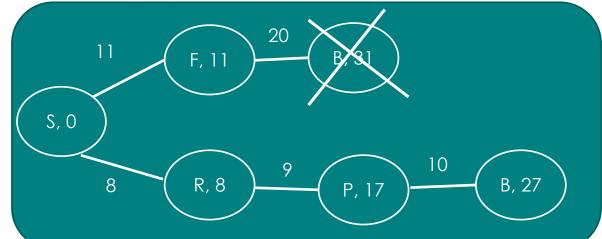
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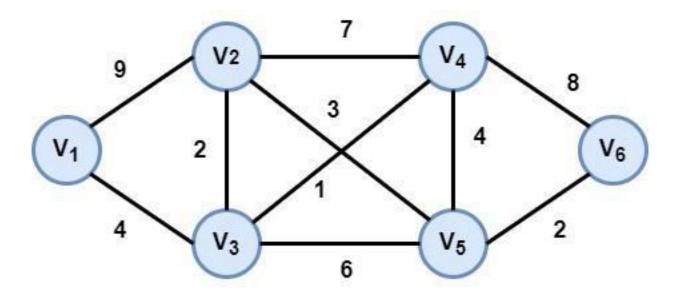
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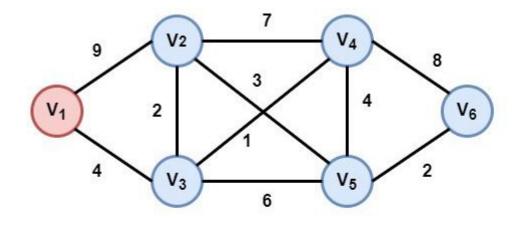
Example

Our target is to go from the city (node) V1 to V6 following the path with the smallest cost (shortest path). Let's execute the UCS algorithm:



☐ Step 1: Initialization

The first node V1 (initial state) of the graph is appended to the opened list. The distance of this node from itself is zero.

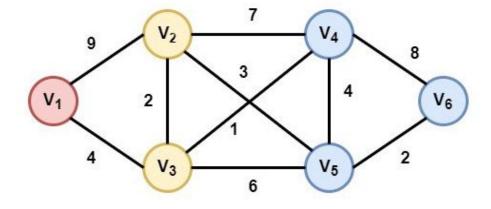


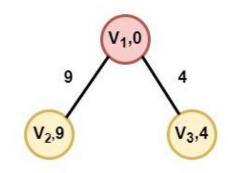
V		1	1
	-	_	

Step	Opened	Explored
0	$\{(V_1,0)\}$	{-}

Step 2: Node V1 is selected

The V1 is selected as it is the only node in the opened list. Its children V2 and V3 are appended in the opened list after the distance calculation from node V1.

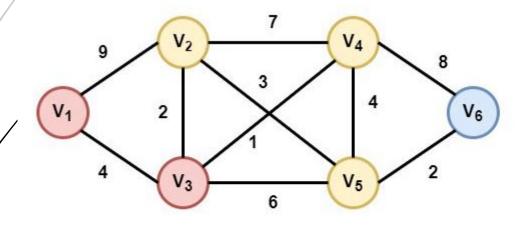




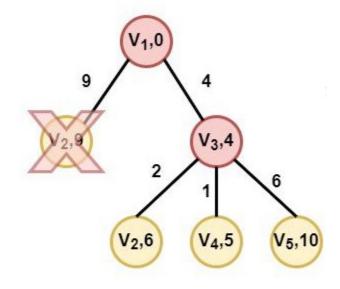
Step	Opened	Explored	
0	$\{(V_1,0)\}$	{-}	
1	$\{(V_2,9),(V_3,4)\}$	$\{(V_1,0)\}$	

Step 3: Node V3 is selected

Node V3 is selected as it has the smallest distance value. As we can see, extending the node V3 we find the node V2 with a smaller distance value. So we replace node $(\sqrt{2},9)$ with the new node $(\sqrt{2},6)$.

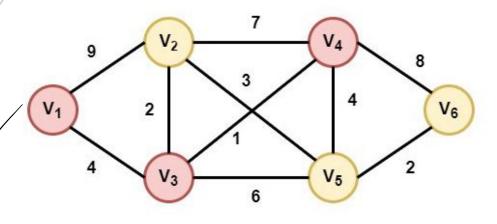


Step	Opened	Explored	
0	$\{(V_1,0)\}$	{-}	
1	$\{(V_2,9),(V_3,4)\}$	$\{(V_1,0)\}$	
2	$\{(V_2,6),(V_4,5),(V_5,10)\}$	$\{(V_1,0),(V_3,4)\}$	

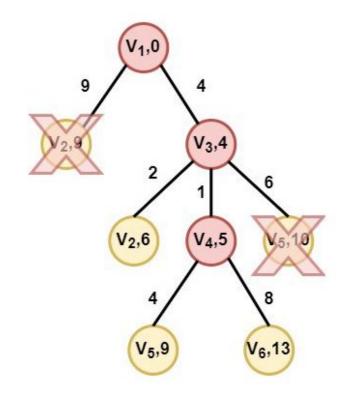


Step 4: Node V4 is selected

Node V4 is selected as it has the smallest distance value. In this step, we find a better distance value for node V5, so we replace the node (V5, 10) with node (V5, 9).

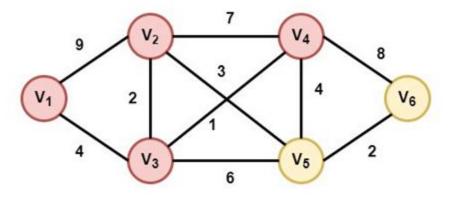


Step	Opened	Explored
0	$\{(V_1,0)\}$	{-}
1	$\{(V_2,9),(V_3,4)\}$	$\{(V_1,0)\}$
2	$\{(V_2,6),(V_4,5),(V_5,10)\}$	$\{(V_1,0),(V_3,4)\}$
3	$\{(V_2,6),(V_6,13),(V_5,9)\}$	$\{(V_1,0),(V_3,4),(V_4,5)\}$

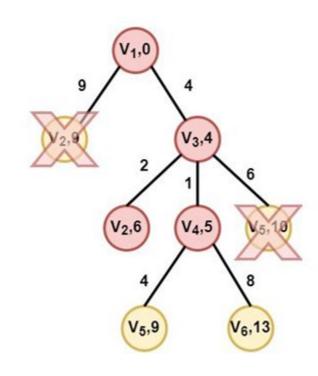


Step 5: Node V2 is selected

Node V2 is selected as it has the smallest distance value. However, none of its children is appended in the opened list, as nodes V3 and V4 are already inserted in the explored list and the algorithm doesn't find a better distance value for node V5.

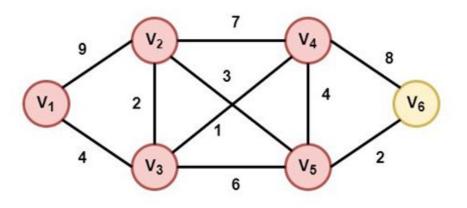


Step	Opened	Explored
0	$\{(V_1,0)\}$	{-}
1	$\{(V_2,9),(V_3,4)\}$	$\{(V_1,0)\}$
2	$\{(V_2,6),(V_4,5),(V_5,10)\}$	$\{(V_1,0),(V_3,4)\}$
3	$\{(V_2,6),(V_6,13),(V_5,9)\}$	$\{(V_1,0),(V_3,4),(V_4,5)\}$
4	$\{(V_6,13),(V_5,9)\}$	$\{(V_1,0),(V_3,4),(V_4,5),(V_2,6)\}$

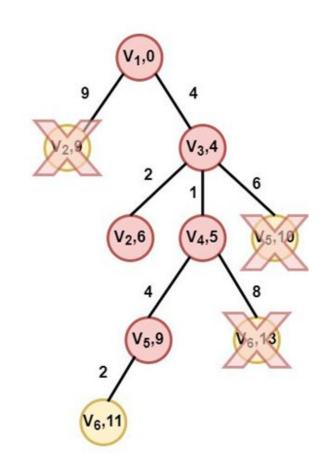


Step 6: Node V5 is selected

Node V5 is selected as it has the smallest distance value. A better path to node V6 is found in this step. So, we replace the old node (V6, 13) with node (V6, 11)

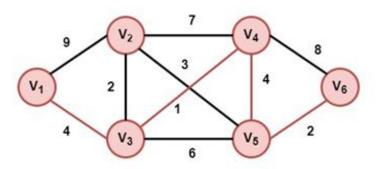


Step	Opened	Explored
0	$\{(V_1,0)\}$	{-}
1	$\{(V_2,9),(V_3,4)\}$	{(V ₁ ,0)}
2	$\{(V_2,6),(V_4,5),(V_5,10)\}$	$\{(V_1,0),(V_3,4)\}$
3	$\{(V_2,6),(V_6,13),(V_5,9)\}$	$\{(V_1,0),(V_3,4),(V_4,5)\}$
4	{(V6,13),(V5,9)}	$\{(V_1,0),(V_3,4),(V_4,5),(V_2,6)\}$
5	{(V ₆ ,11)}	$\{(V_1,0),(V_3,4),(V_4,5),(V_2,6),(V_5,9)\}$

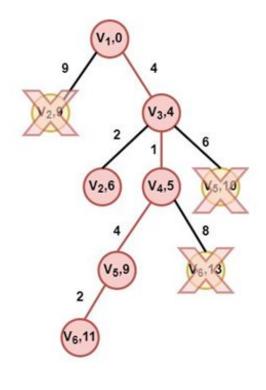


Step 7: Node V6 is selected

Node V6 (target node) is selected. So the algorithm returns the path from node V1 to node V6 with cost 11, which constitutes the best solution.



Step	Opened	Explored
0	$\{(V_1,0)\}$	{-}
1	$\{(V_2,9),(V_3,4)\}$	{(V ₁ ,0)}
2	$\{(V_2,6),(V_4,5),(V_5,10)\}$	$\{(V_1,0),(V_3,4)\}$
3	$\{(V_2,6),(V_6,13),(V_5,9)\}$	$\{(V_1,0),(V_3,4),(V_4,5)\}$
4	$\{(V_6,13),(V_5,9)\}$	$\{(V_1,0),(V_3,4),(V_4,5),(V_2,6)\}$
5	{(V ₆ ,11)}	$\{(V_1,0),(V_3,4),(V_4,5),(V_2,6),(V_5,9)\}$
6	{-}	$\{(V_1,0),(V_3,4),(V_4,5),(V_2,6),(V_5,9)\}$



Path: V1 - V3 - V4 - V5 - V6

Total Cost: 11

When to use BFS? Real-Life Applications

Crawlers in Search Engines:

BFS is used for indexing web pages. The algorithm starts traversing from the **source page and follows all the links associated with the page**. Here each web page will be considered as a node in a graph.

GPS Navigation systems:

BFS is used to find **neighboring locations** by using the GPS system.

Broadcasting:

Networking used to communicate **broadcasted packets** across all the nodes in a network. These packets follow a traversal method of BFS to reach various networking nodes.

Peer to Peer Networking:

BFS can be used as a traversal method to find all the **neighboring nodes** in a Peer to Peer Network. For example, **BitTorrent uses Breadth-First Search for peer to peer communication**.

When to use DFS? Real-Life Applications

Cycle detection:

DFS can be used to detect cycles in a graph. If a node is visited again during a DFS traversal, it indicates **that there is a cycle in the graph**.

Pathfinding:

DFS can be used to find a path between two nodes in a graph.

Solving puzzles:

DFS can be used to **solve puzzles such as mazes**, where the goal is to find a path from the start to the end.

Backtracking:

DFS can be used for **backtracking in algorithms like Sudoku**.

