

Advanced Artificial Intelligence

Week #3

Dr. Qurat Ul Ain

Assistant Professor Dept. of Al & DS FAST NUCES, Islamabad Email:

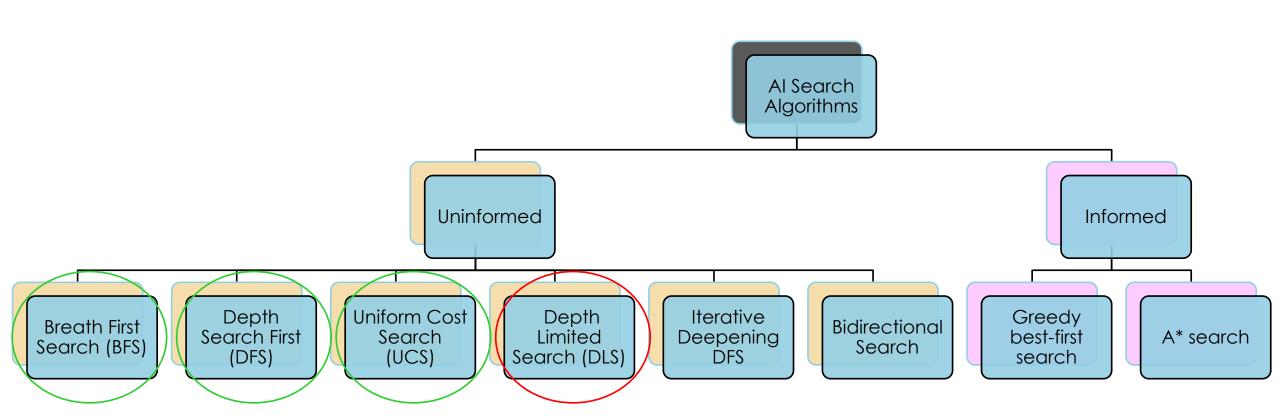




Uninformed Search Algorithms in Artificial Intelligence

Learning Objective of this Topic

- Uninformed Searching Algorithms
 - Uniform Cost Search (UCS)
 - Depth Limited Search (DLS)
 - ☐ Iterative Deepening Search (IDS)
 - ☐ Bidirectional Search (BS)



Depth Limited

Search (DLS)

What is Depth-Limited Search (DLS)?

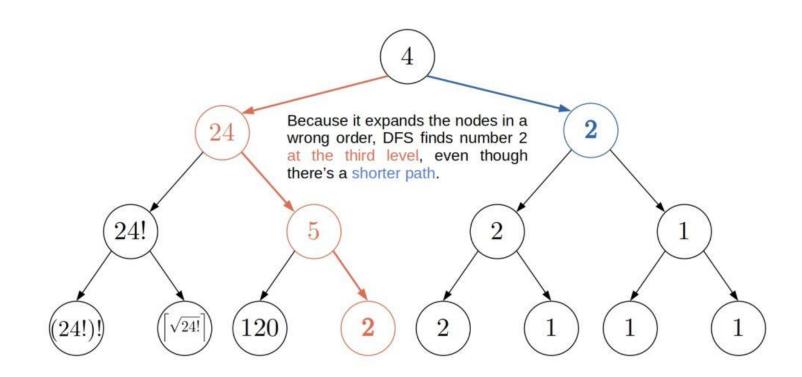
☐ The depth-limited search is a variation of a well-known depth-first search(DFS) traversing algorithm.

☐ It performs depth-first search (DFS) by implementing a depth limit.

But why we need Depth Limited Search?

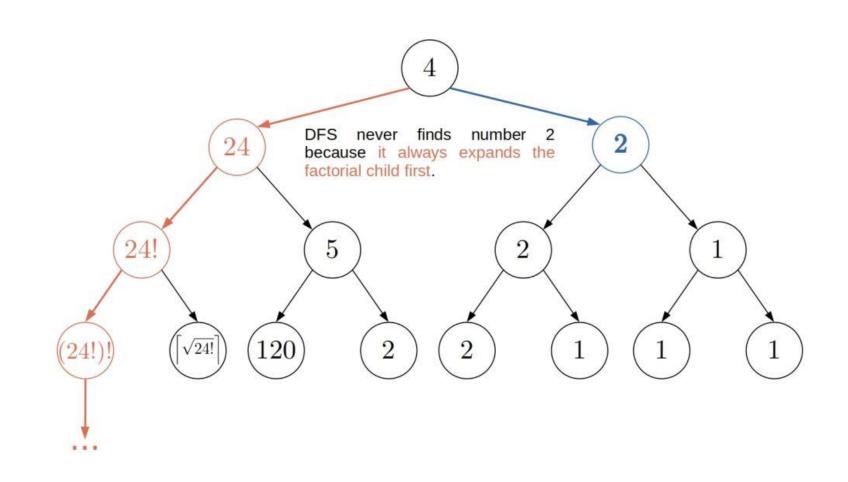
1- Problem with Depth First Search

DFS may miss the optimal path. Depending on the order in which children returns the children of a node, DFS may expand more nodes than necessary:



2- Problem with Depth First Search

Also, DFS may never end! It may get stuck at expanding the nodes that can't lead to a target node



Depth Limited Search

- limit: the depth of the tree/graph until which we explore the nodes
- Expand Depth Node First (Depth Limit Search)
- Stack: nodes in the stack to be explored
- explored: Nodes that are already explored

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function DEPTH-LIMITED-SEARCH(problem, limit) returns a node or failure

Stack ← a LIFO with source node as the only element

explored ← an empty set

loop do

if EMPTY?(Stack) then return failure

node ← POP(Stack) /* chooses the top node in Stack*/

add node to explored

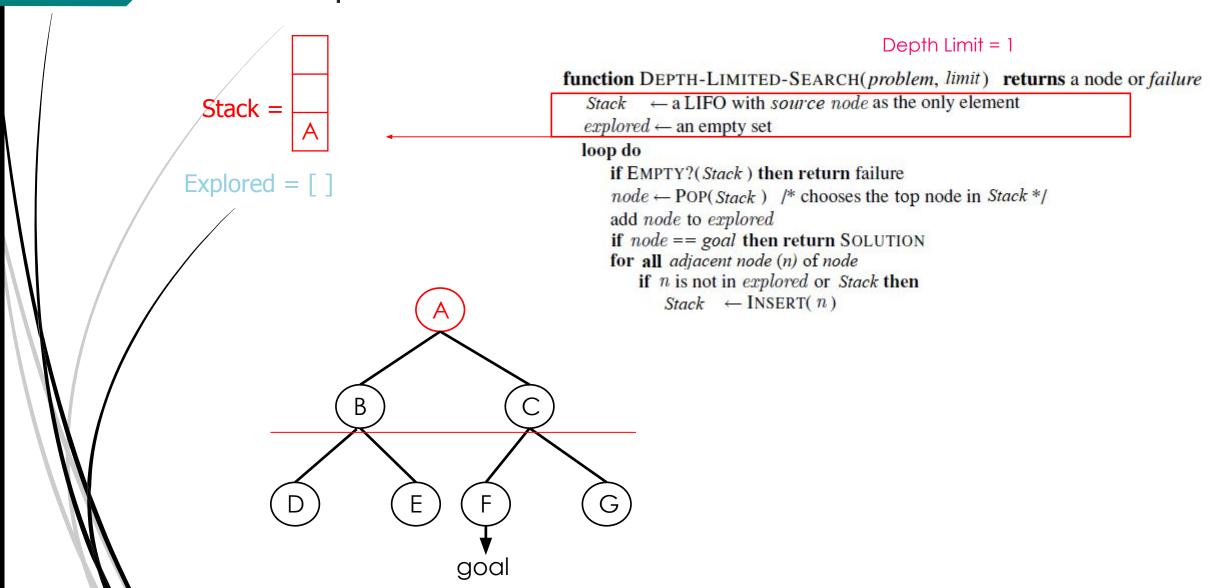
if node == goal then return SOLUTION

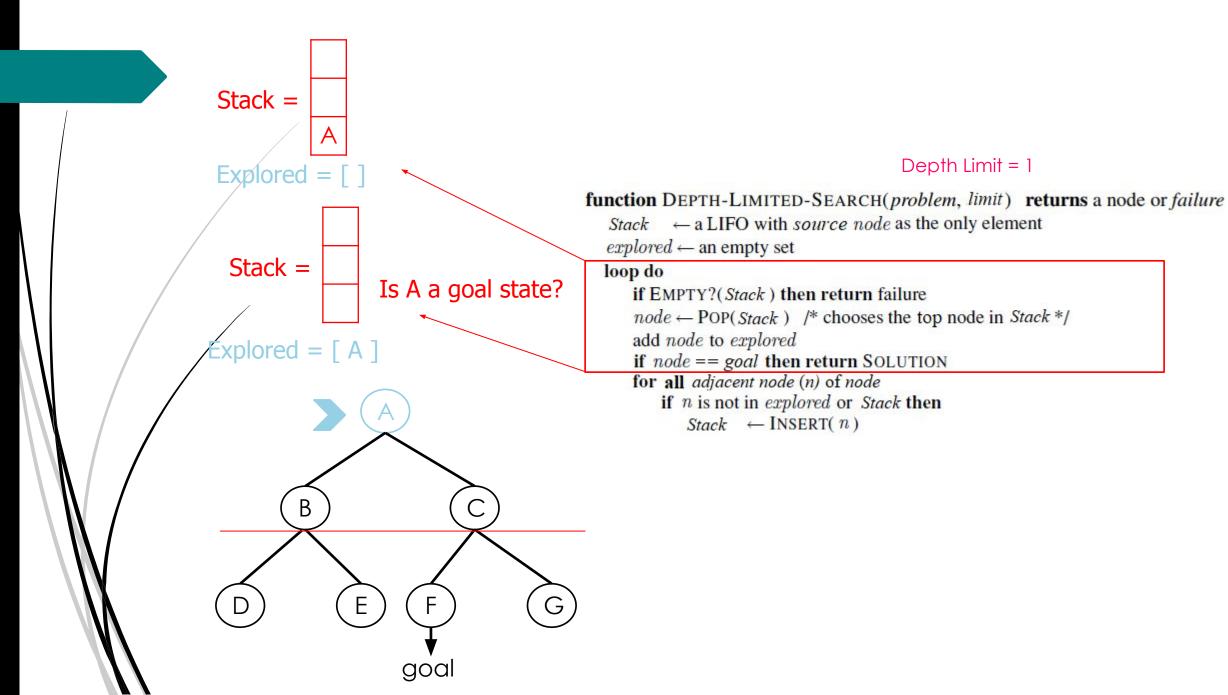
for all adjacent node (n) of node

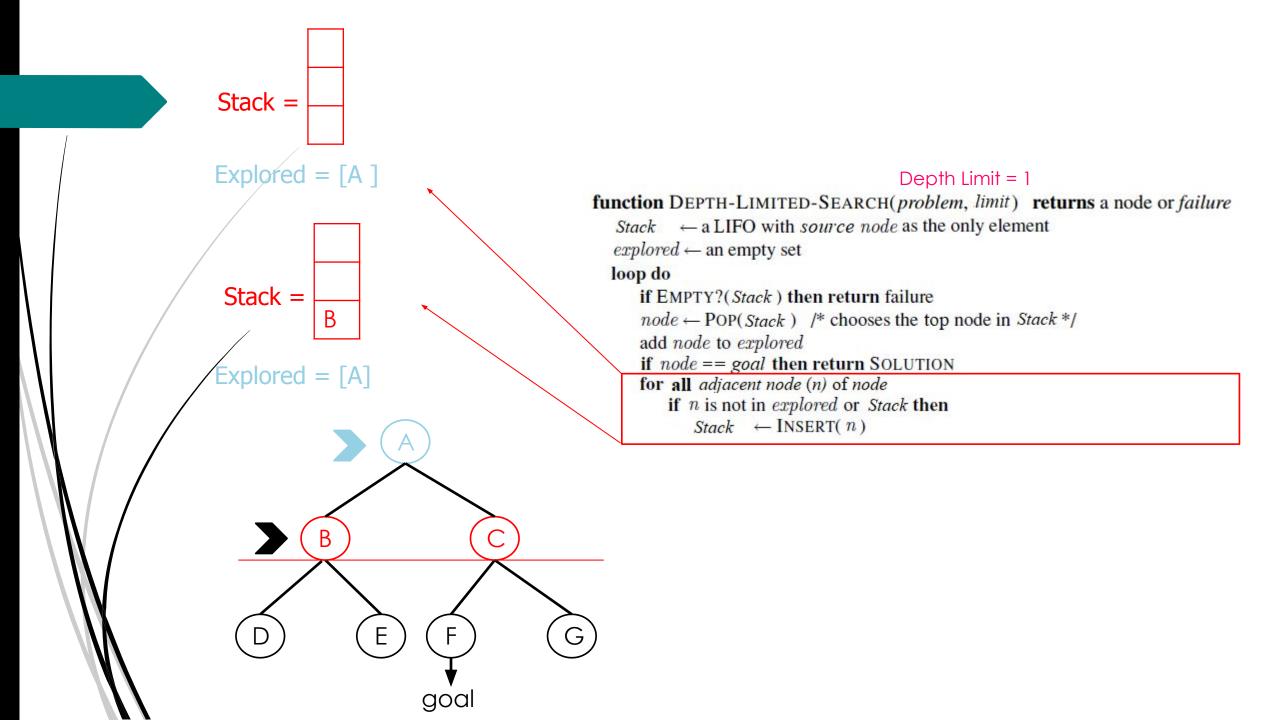
if n is not in explored or Stack then

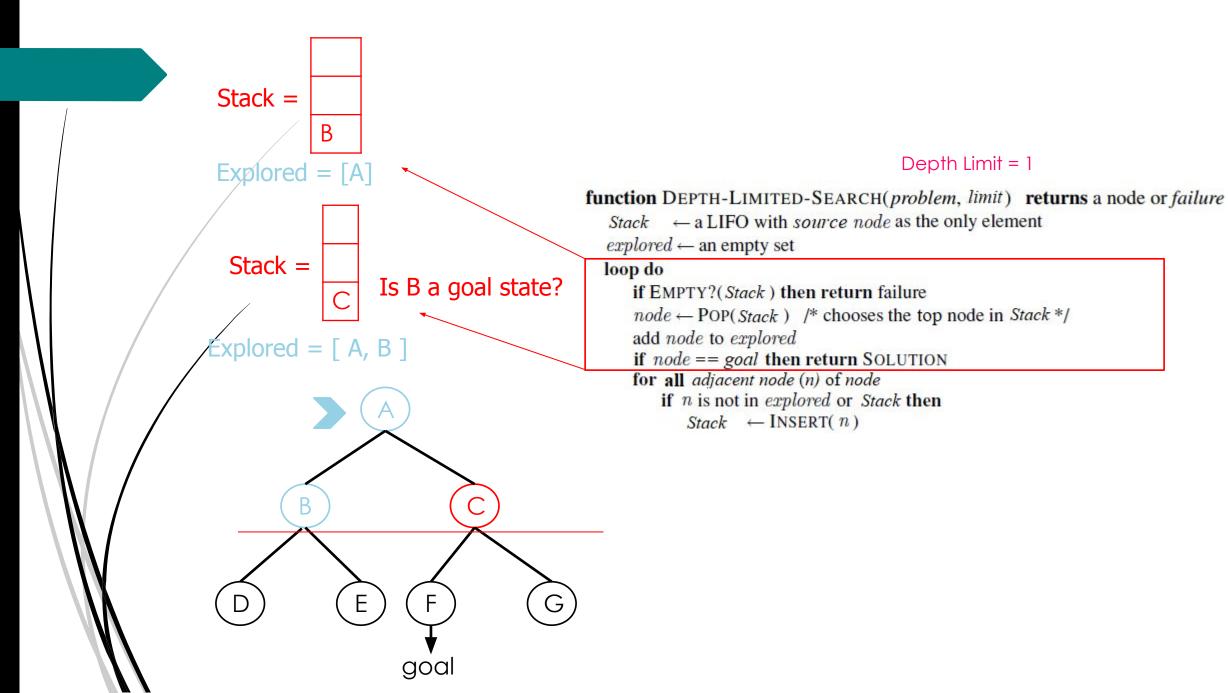
Stack ← INSERT(n)
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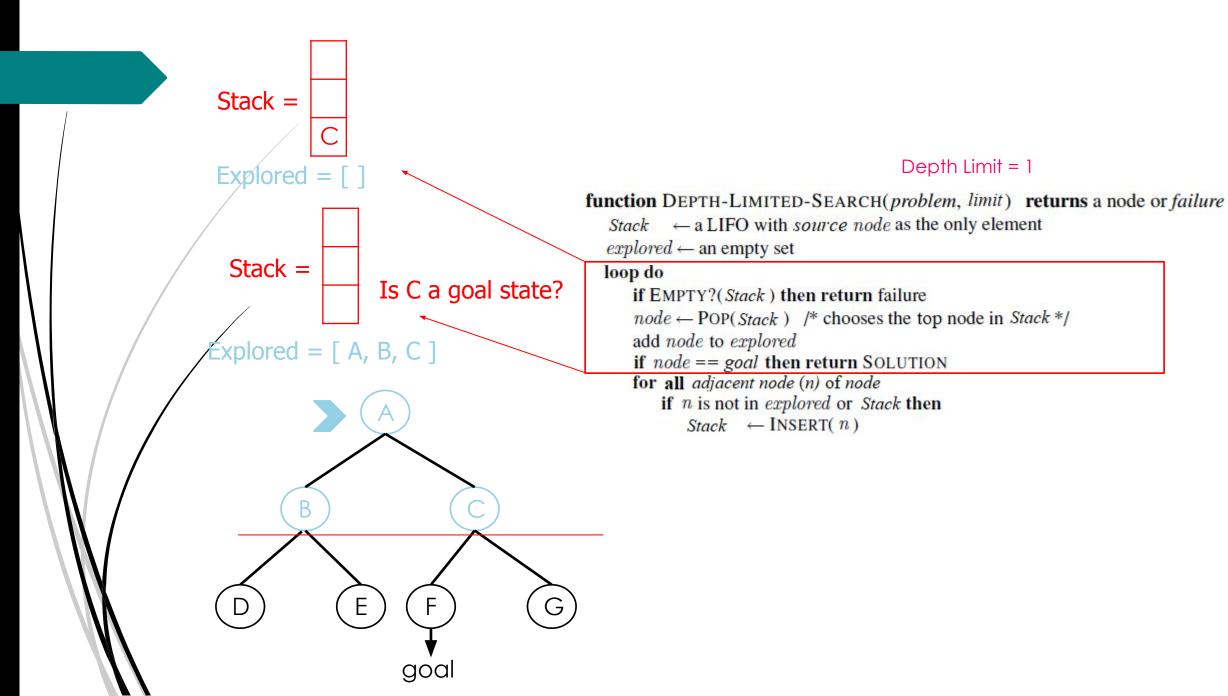
Depth Limited Search

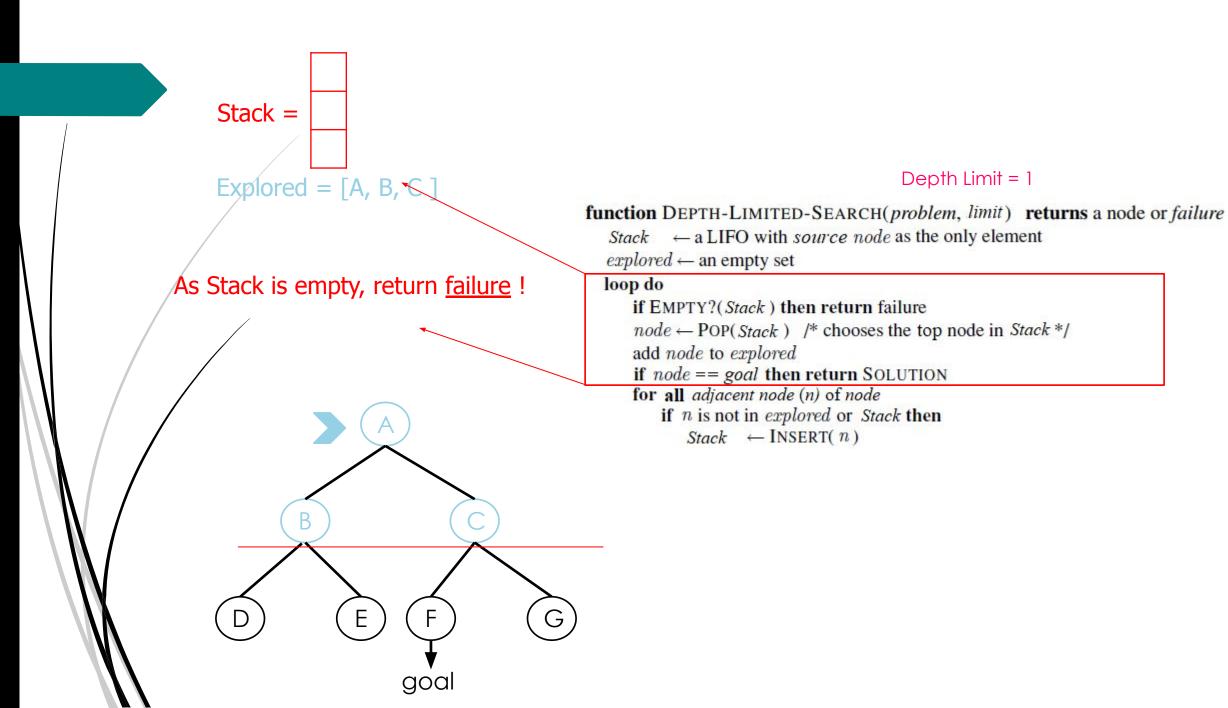


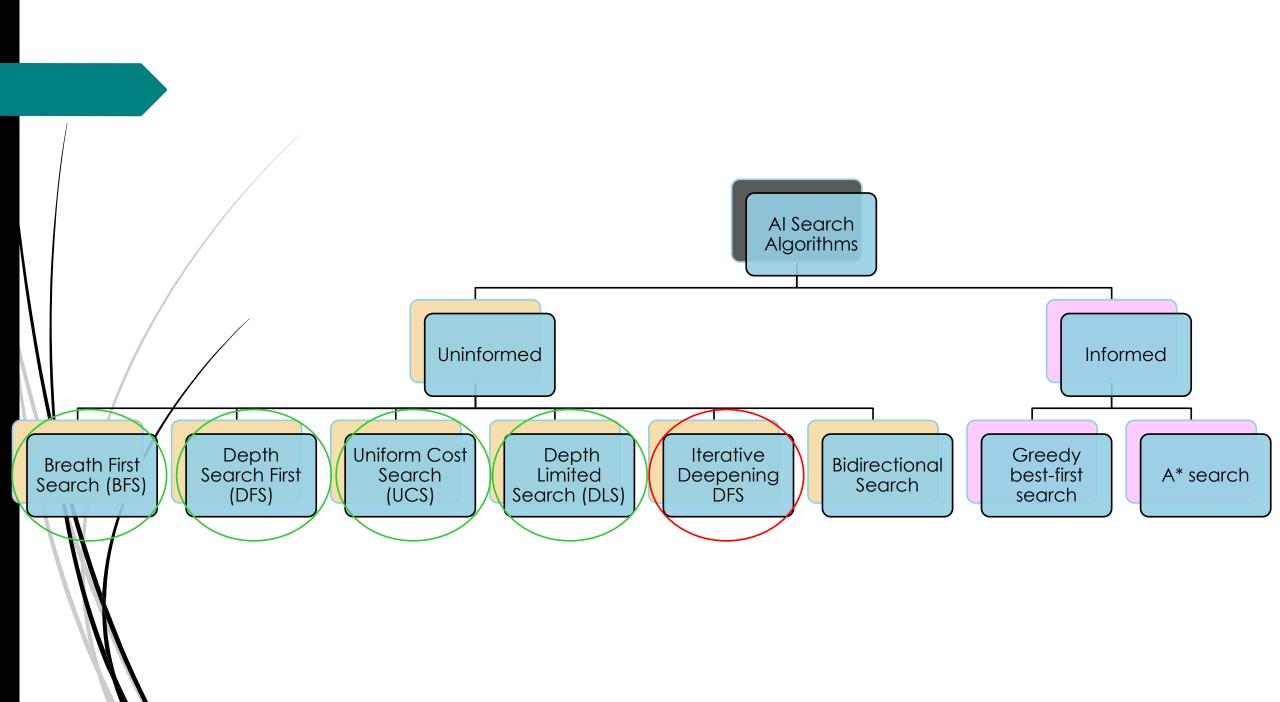


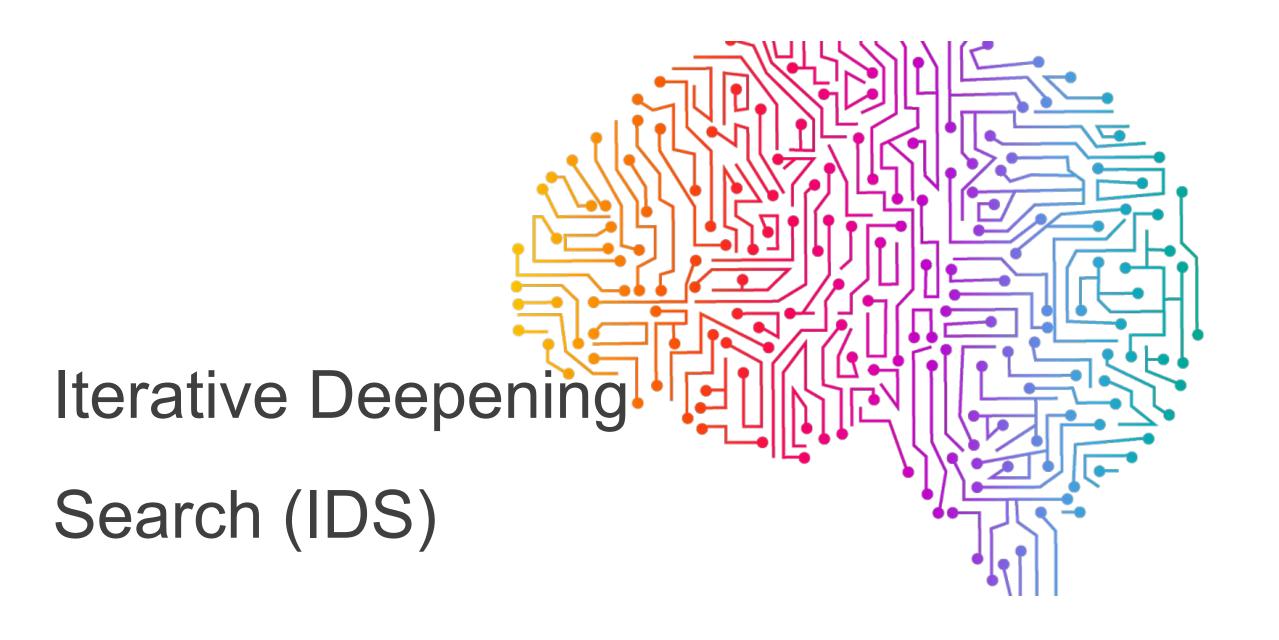












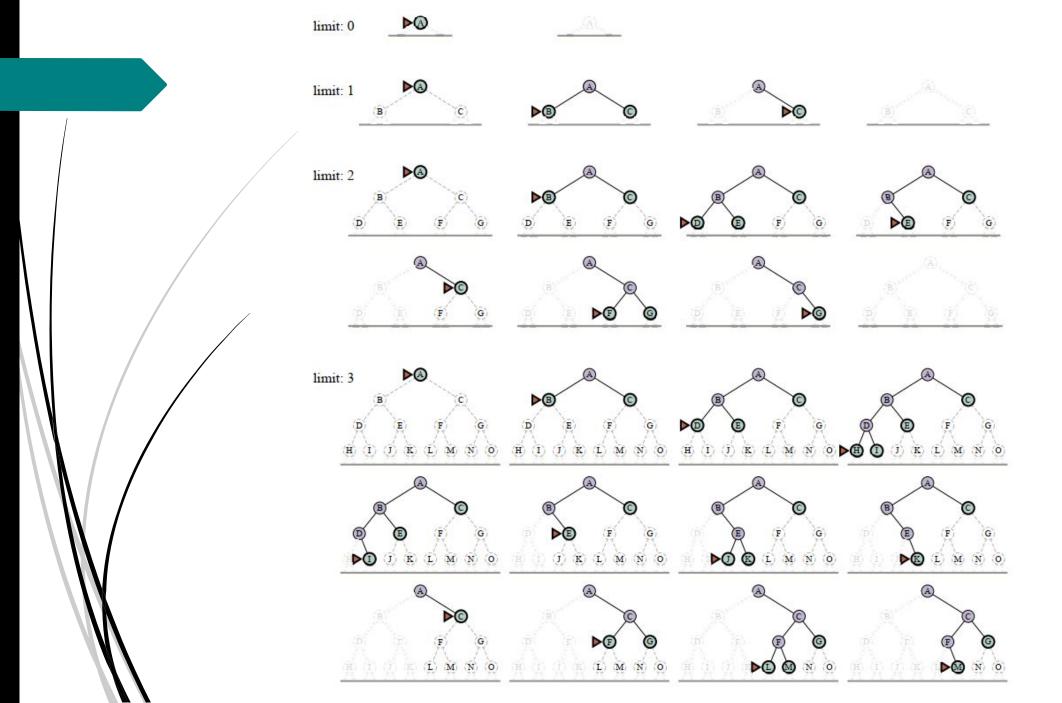
Iterative Deepening Search

- What if the solution is deeper than Limit?
 - Increase depth iteratively
 - Iterative Deepening Search
- Iterative deepening search solves the problem of picking a good value for depth limit by trying all values: first 0, then 1, then 2, and so on—until either a solution is found, or the depth-limited search returns the failure.
- Also known as Iterative Deepening Depth First Search or Iterative
 Deepening Depth Limited Search.

Iterative Deepening Search

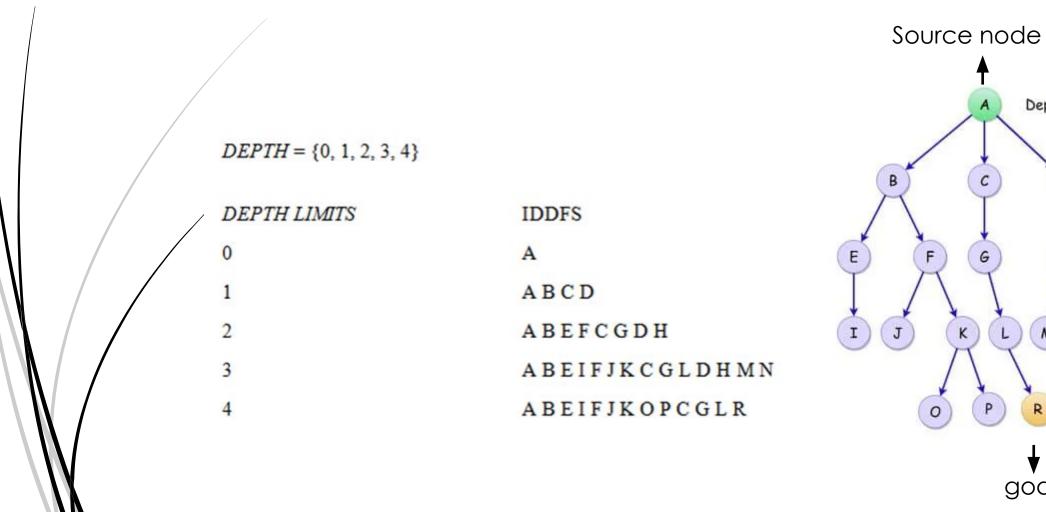
- Also known as Iterative Deepening Depth First Search or Iterative Deepening Depth Limited Search.
- IDS GENERALLY THE PREFERRED UNINFORMED SEARCH
 - ☐ Inherits the memory advantage of depth-first search
 - Has the completeness property of breadth-first search

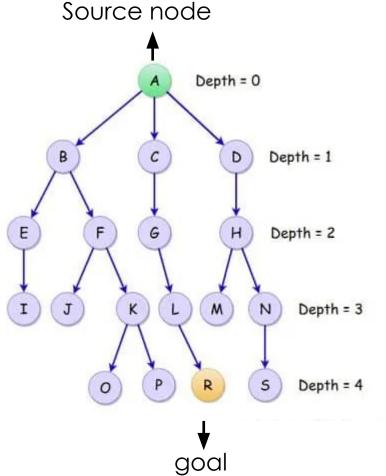
function ITERATIVE-DEEPENING-SEARCH(problem) returns a solution node or failure
for limit = 0 to ∞ do
 result ← DEPTH-LIMITED-SEARCH(problem, limit)

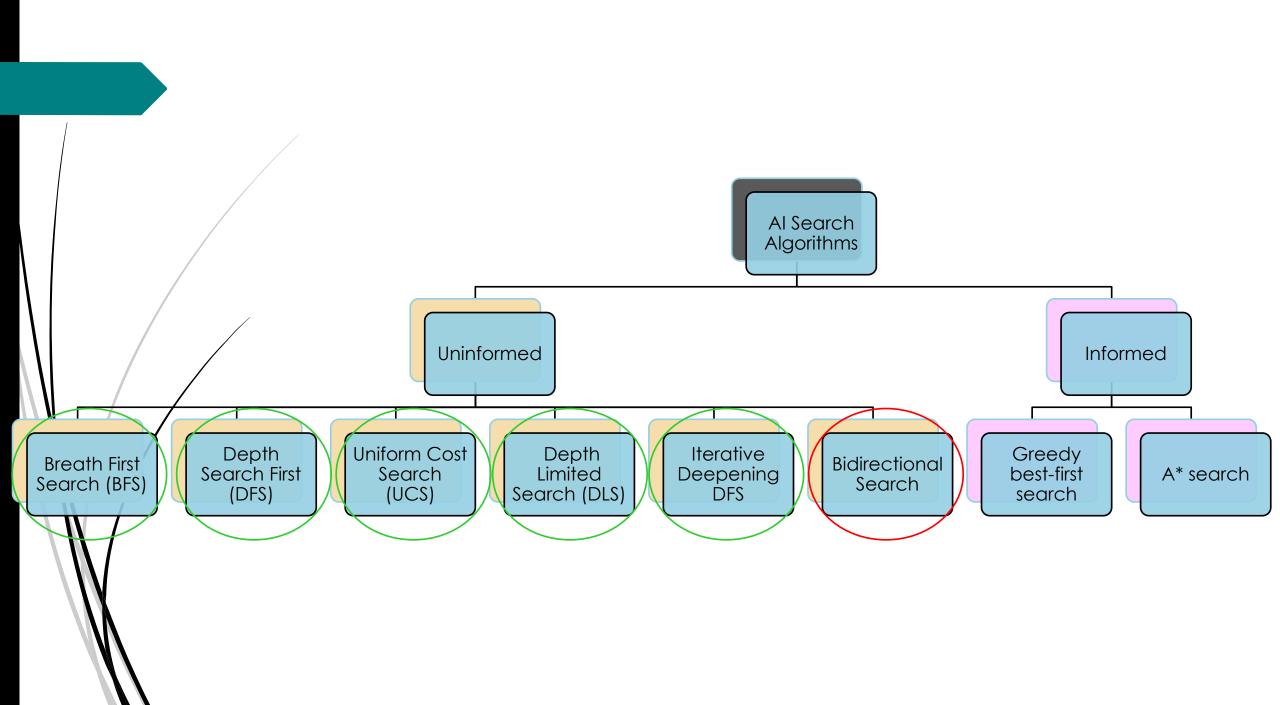


M is the goal node

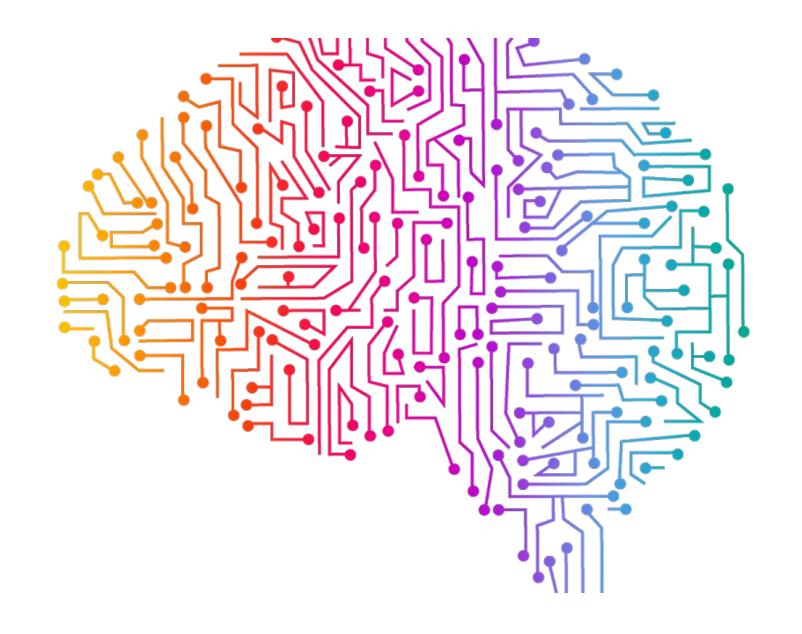
Example IDS







Bidirectional
Search (BS)



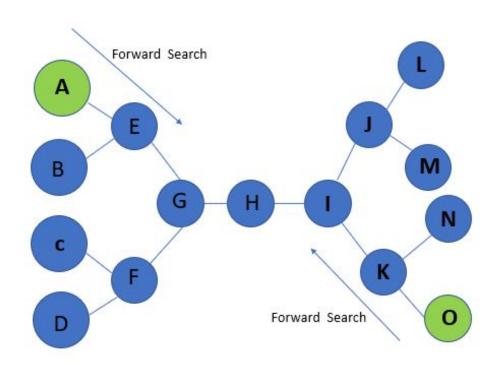
Bidirectional Search

Bidirectional search is a graph search algorithm that finds the smallest path from the source to the goal vertex. It runs two simultaneous searches –

- ☐ Forward search from source/initial vertex toward goal vertex
- Backward search from goal/target vertex toward source vertex

Steps in Bidirectional Search

- Step 1: Say, A is the initial node and O is the goal node, and H is the intersection node.
- Step 2: We will start searching simultaneously from start to goal node and backward from goal to start node.
- Step 3: Whenever the forward search and backward search intersect at one node, then the searching stops.



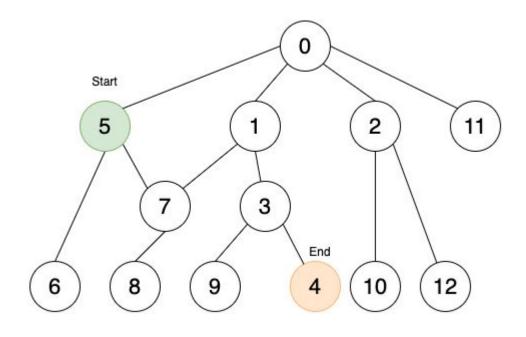
Completeness: Bidirectional search is complete if BFS is used in both searches.

Optimality: It is optimal if BFS is used for search and paths have uniform cost.

Example

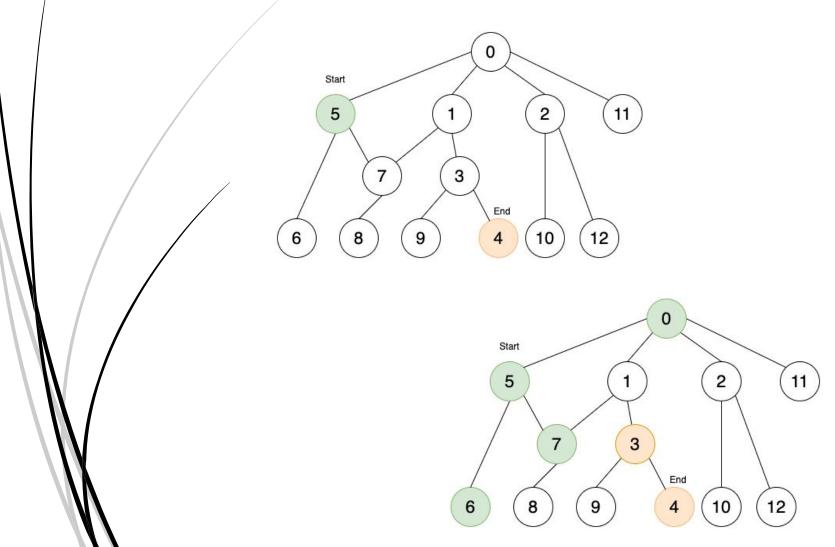
The start node is 5 and the end node is 4.

Aim: To find the shortest path from 5 to 4 using bidirectional search.

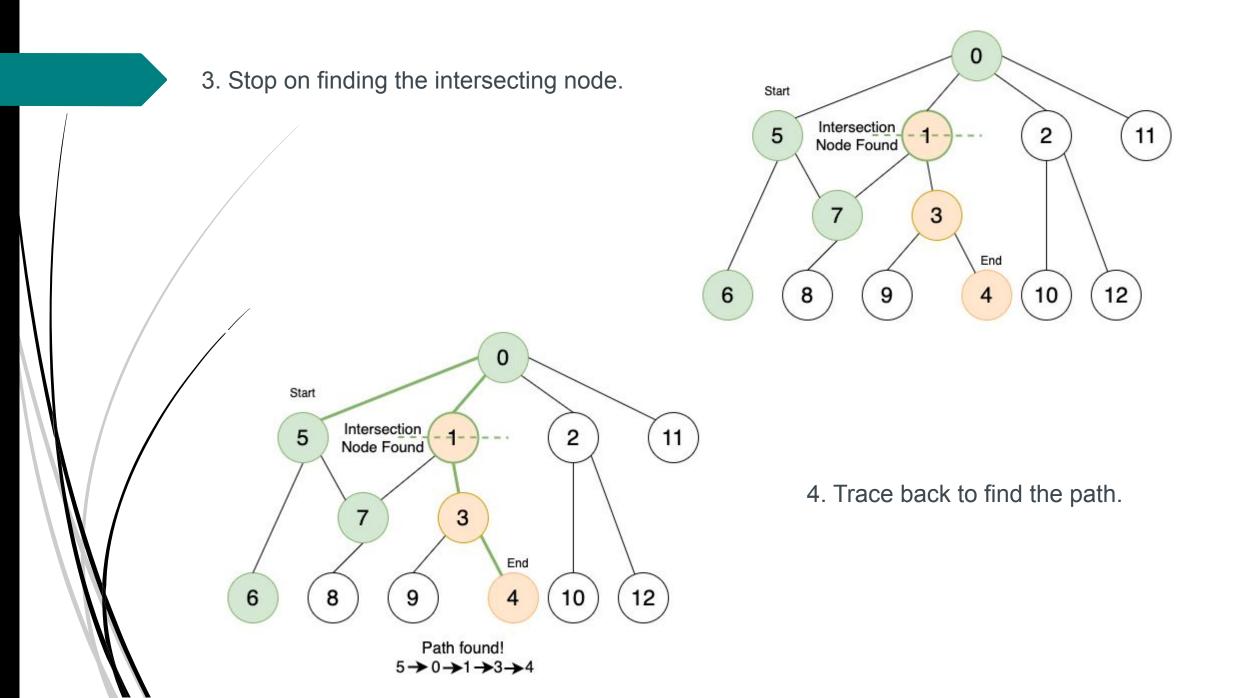


Do BFS from both directions.

1.Start moving forward from start node (Green) and backwards from end node (Orange).



2. Similar to BFS, at every point explore the next level of nodes till you find an intersecting node.



Comparing
Uninformed

Searching Algorithms

Comparing Uninformed Searching Algorithms

Criterion	Breadth- First	Uniform- Cost	Depth- First	Depth- Limited	Iterative Deepening	Bidirectional (if applicable)
Complete? Optimal cost? Time Space	Yes^1 Yes^3 $O(b^d)$ $O(b^d)$	$\operatorname{Yes}^{1,2} \ \operatorname{Yes} \ O(b^{1+\lfloor C /\epsilon \rfloor}) \ O(b^{1+\lfloor C /\epsilon \rfloor})$	No No $O(b^m)$ $O(bm)$	No No $O(b^\ell)$ $O(b\ell)$	Yes^1 Yes^3 $O(b^d)$ $O(bd)$	${ m Yes}^{1,4} \ { m Yes}^{3,4} \ O(b^{d/2}) \ O(b^{d/2})$

Evaluation of search algorithms. b is the branching factor; m is the maximum depth of the search tree; d is the depth of the shallowest solution, ℓ is the depth limit.

¹ complete if b is finite, and the state space either has a solution

² complete if all action costs are $\geq \epsilon > 0$;

³ cost-optimal if action costs are all identical:

⁴ if both directions are breadth-first or uniform-cost.

 $[\]epsilon$ is an arbitrarily small positive quantity

