

Tim Marder

APCS2 Pd02

HW#06 -- How Fast Are Your Turtles?

2018-02-13

Win

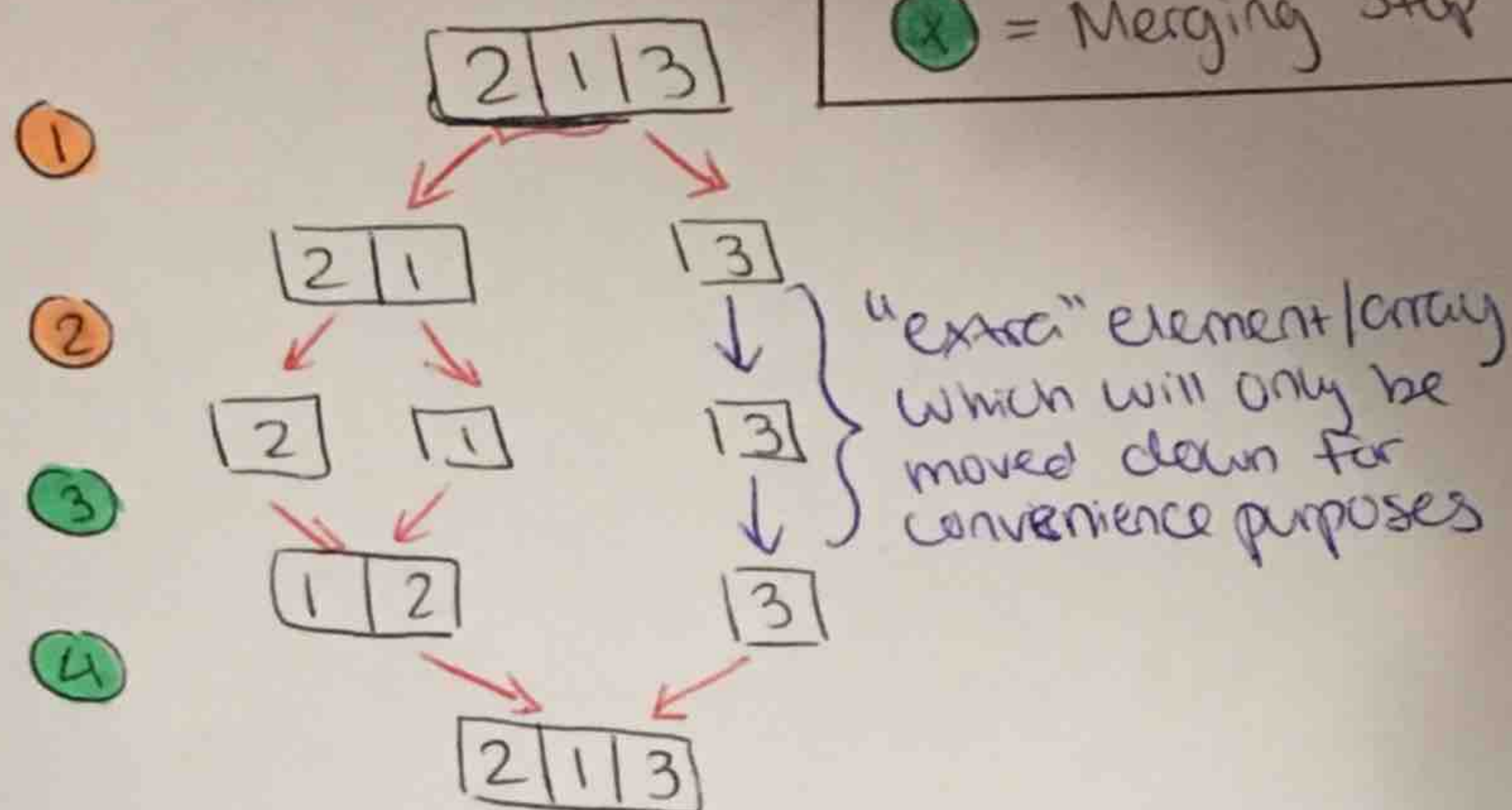
Key

→ = Split or Merge

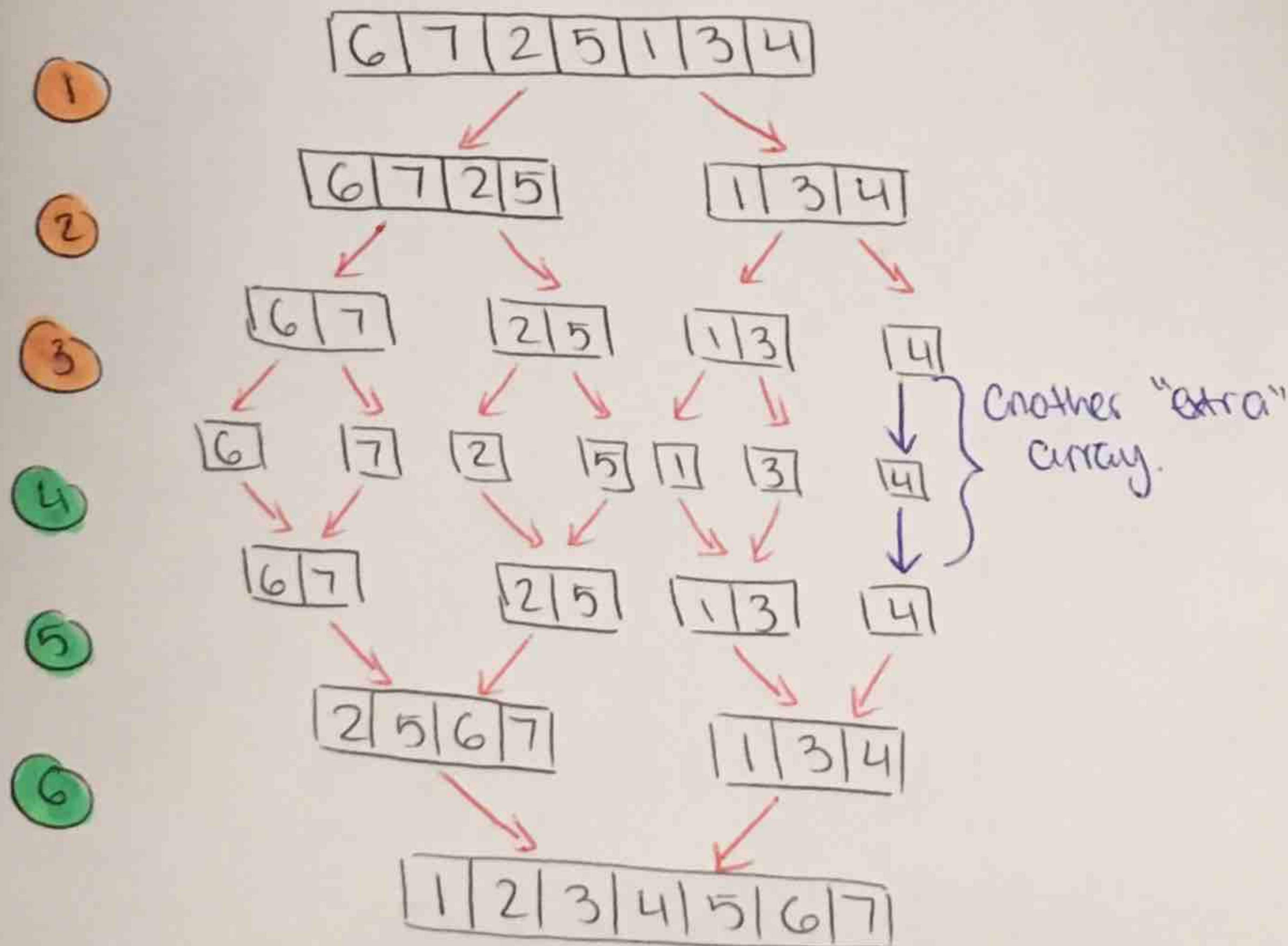
→ = No action; just brought down elements (arrays)

⊗ = Splitting Step

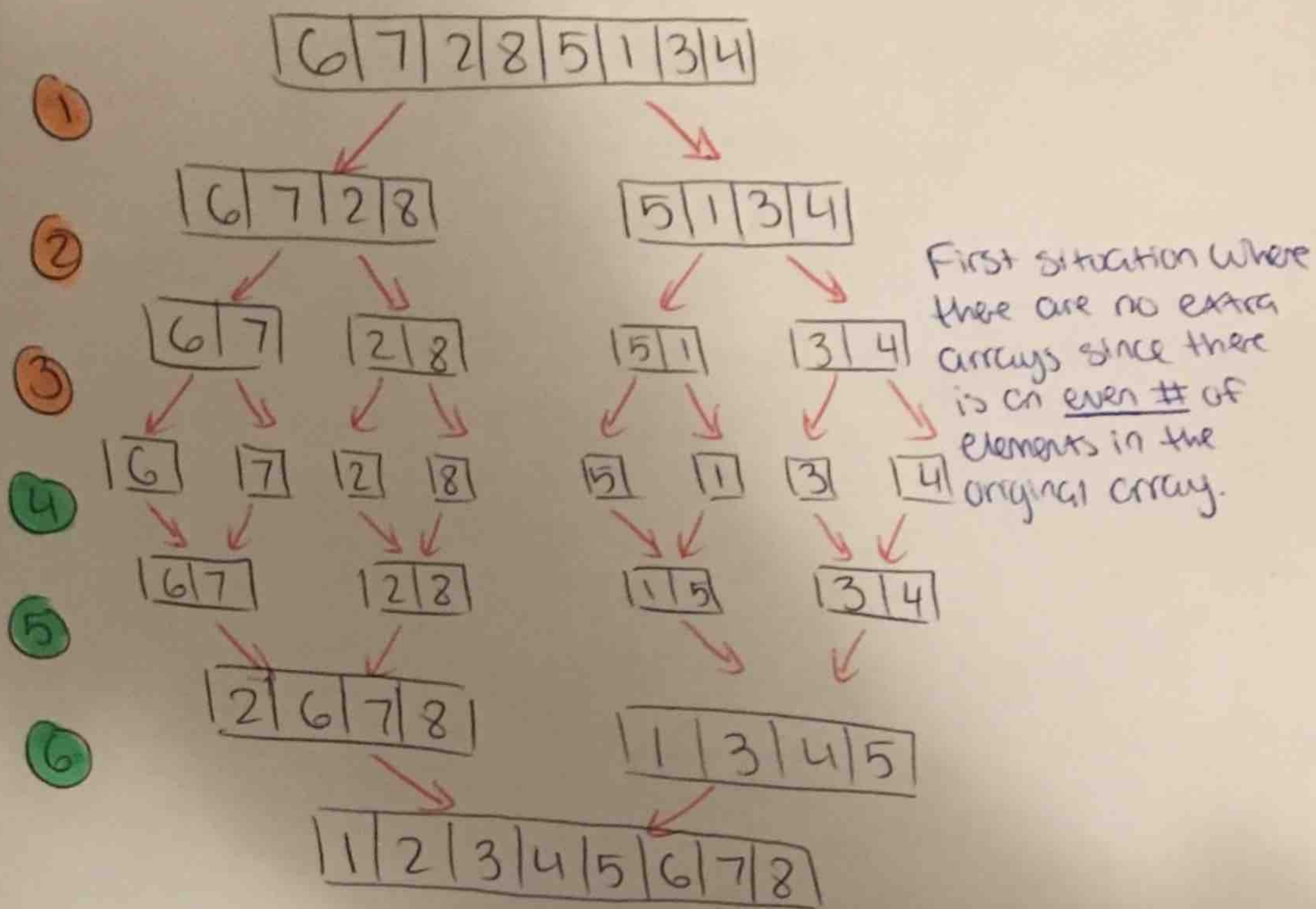
⊙ = Merging Step



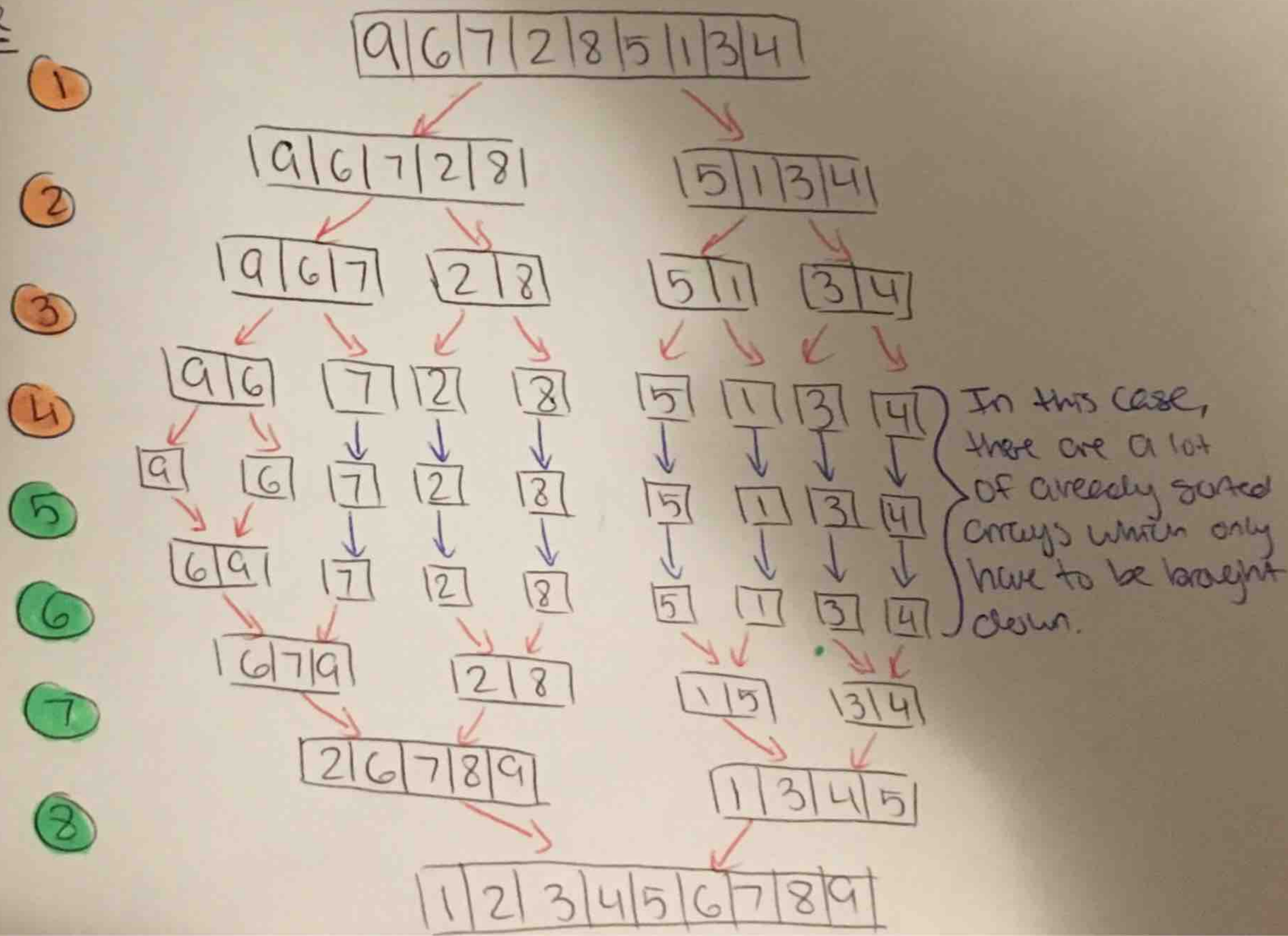
Tim :



Jeremy



he'gine



Richard

0 | 6 | 7 | 2 | 8 | 5 | 1 | 9 | 3 | 4

0 | 6 | 7 | 2 | 8

First situation
where original
array is
split into 2
arrays
with an
odd #
of
elements

5 | 1 | 9 | 3 | 4

0 | 6 | 7

2 | 8

5 | 1 | 9

3 | 4

0 | 6

7

2

8

5 | 1

9

3

4

0

6

7

2

8

5

1

9

3

4

0 | 6

7

2

8

1 | 5

9

3

4

0 | 6 | 7

2 | 8

1 | 5 | 9

3 | 4

0 | 2 | 6 | 7 | 8

1 | 3 | 4 | 5 | 9

0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9

1

2

3

4

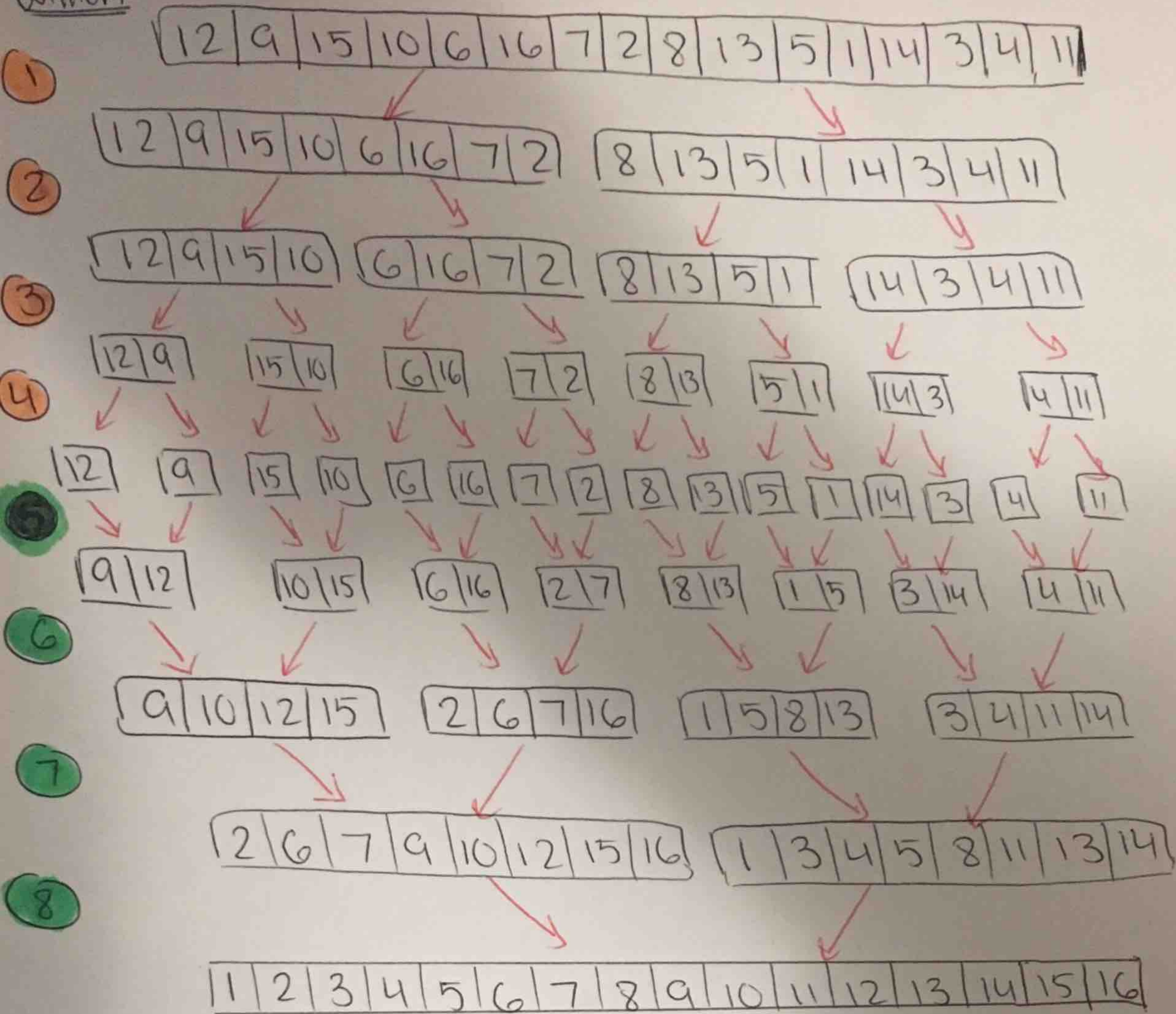
5

6

7

8

William:



The execution time of this merge sort algorithm is $O(n * \log_2 n)$. This can be seen by some of the arrays above that had to be sorted using merge sort. The amount of steps required to completely sort an array would be equal to $2^x = n$ (n being # of elements in original array). The value of 2^x must be greater or equal to n to be enough for the array to fully sort. The largest example (William) has 16 elements, which is 2^4 , but if one more element was added, it would have to become 2^5 , and it would stay 2^5 until the # of elements would exceed 32. Arrays "Tim" and "Jeremy" are 2^3 since the # of elements is larger than 4 but less than or equal to 8. So the splitting part of the algorithm is $O(\log_2 n)$ since the array gets divided into HALVES. The merging part of the algorithm is just $O(n)$ since it is linear, like discussed earlier. So the resulting execution time is $O(n * \log_2 n)$.