

# Simulating the Heisenberg Model

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We estimate the entanglement entropy of the GHZ state using both simulated and real measurement data. We use Qiskit to measure the target state in the Pauli basis and then use the classical shadows protocol to estimate the entanglement entropy from our measurement data. Our estimated value aligns very well with theory for our simulated data (relative error  $\leq 0.15\%$ ), but our estimated value based on real data differs substantially from the theoretical value for our target state (relative error  $\approx 20\%$ ). We conclude that the device's large error is most likely caused by noise.

## INTRODUCTION

### EXPERIMENTAL HARDWARE

For our simulation task we will be using the `ibmq-jakarta` device. This is a 7 qubit quantum computer whose qubits are connected as depicted in figure 1. Our simulation task will require us to use gates on this device, but like many other devices `ibmq-jakarta` has a limited gate set. This means that any gate we want to execute on the device must either be in the gate set or must be broken down into some composition of gates in the gate set. The gate set on this device consists of  $R_z(\theta)$ ,  $X$ ,  $\sqrt{X}$ , and  $CNOT$  gates. The first gate  $R_z(\theta)$  is an arbitrary single qubit rotation about the z-axis. The second is the  $\sigma^x$  or NOT gate. The third  $\sqrt{X}$  is the square root of the NOT gate, and the last  $CNOT$  is the controlled not gate. We note that  $\sqrt{X}$  is equivalent (differs only by a global phase) to a  $\pi/2$  rotation about the x-axis  $R_x(\pi/2)$ .

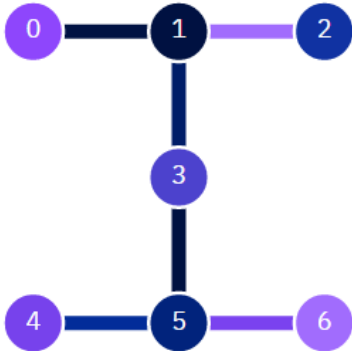


FIG. 1. This figure depicts how the qubits of the `ibmq-jakarta` device are connected to one another. Two qubit gates can only be performed on connected qubits.

## QUANTUM SIMULATION

The Heisenberg model, in its most general form, is defined as

$$H = - \sum_{j=1}^N (J_x \sigma_j^x \sigma_{j+1}^x + J_y \sigma_j^y \sigma_{j+1}^y + J_z \sigma_j^z \sigma_{j+1}^z + h \sigma_j^z), \quad (1)$$

where  $h$  is the external magnetic field and  $J_x$ ,  $J_y$ , and  $J_z$  are the couplings along the x, y, and z directions respectively. The subscript  $j$  denotes the site that each spin operator is acting on, and  $\sigma$ 's are the pauli matrices which are given by

$$\sigma^x = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}, \sigma^y = \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix}, \sigma^z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix} \quad (2)$$

when written in the z-basis. We will be considering the XXX Heisenberg model where  $J = J_x = J_y = J_z$  and  $h = 0$ . For simplicity, we will assume  $J = 1$  throughout this paper.

The Heisenberg model is useful for studying the dynamics and behavior of quantum mechanical magnetic interactions of spins. We will be simulating the quantum mechanical version of the Heisenberg model on three spin- $\frac{1}{2}$  particles. For our task, we start with the initial condition

$$|\psi(t=0)\rangle = |110\rangle \quad (3)$$

with open boundary conditions and seek to simulate a total amount of time  $\tau = \pi$ . Theoretically, the time evolution of this initial condition is determined by the time-evolution operator

$$|\psi(t=\tau)\rangle = e^{-i\hat{H}\tau} |\psi(0)\rangle, \quad (4)$$

where we have set  $\hbar = 1$ . For our problem the time evolution operator is the exponential of a sum of pauli strings. For instance, we could generically write the hamiltonian as

$$H = \sum_i h_i \quad (5)$$

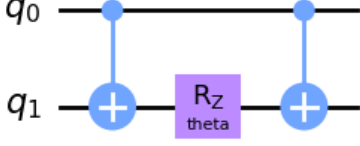


FIG. 2. This circuit shows a z-rotation sandwiched between two CNOT gates. This circuit is equivalent to the two qubit rotation given in equation (6).

where each of the  $h_i$ 's is the tensor product of two pauli matrices on two nearest neighbor sites. It is fairly straightforward to implement the unitary evolution of each  $h_i$  taken individually. For instance, suppose we wanted to implement

$$e^{-i\sigma_1^z \otimes \sigma_2^z t}. \quad (6)$$

Our native gate set lets us do single-qubit z-rotations which are defined as

$$R_z(2t) = e^{-i\sigma_1^z \otimes I t}, \quad (7)$$

where the identity is made explicit for clarity. We can turn this into a two qubit rotation by noticing that a  $Z$  gate with a CNOT on either side is equivalent to a  $Z$  gate on two qubits, i.e.,  $\text{CNOT} (\sigma^z \otimes I) \text{CNOT} = \sigma^z \otimes \sigma^z$ . This relationship still holds when  $\sigma^z \otimes I$  is in an exponential (see figure 2). Additionally, by using the Hadamard gate  $H$  and the phase gate  $S$  we can exploit relationships such as  $HZH = X$  and  $SXS^\dagger = Y$  to build other unitaries generated by individual  $h_i$ . This is not sufficient, however. The different  $h_i$  do not commute, so we cannot split the exponential of the sum into the product of exponentials. It is for this reason that we must use a technique known as trotterization.

Generally speaking, if we have two matrices  $h_1$  and  $h_2$

$$e^{(h_1+h_2)t} \neq e^{h_1 t} e^{h_2 t}. \quad (8)$$

Though the two sides of this equation are not equal, they are approximately true when  $t$  is small. If we want to simulate to a time  $t$ , we need only split the total time into  $N$  steps and apply the two gates successively. More precisely

$$e^{(h_1+h_2)t} = \lim_{n \rightarrow \infty} \left( e^{h_1 t/n} e^{h_2 t/n} \right)^n \quad (9)$$

$$\approx \left( e^{h_1 t/N} e^{h_2 t/N} \right)^N. \quad (10)$$

Using the Zassenhaus formula

$$\begin{aligned} e^{t(h_1+h_2)} &= e^{th_1} e^{th_2} e^{-\frac{t^2}{2}[h_1, h_2]} \dots \\ &= e^{th_1} e^{th_2} \left( 1 - \frac{t^2}{2}[h_1, h_2] + \dots \right) \\ &= e^{th_1} e^{th_2} + \mathcal{O}(t^2), \end{aligned}$$

we can see that this trotterization method has a second-order accuracy [1].

Recalling the commutation relations for pauli matrices

$$[\sigma^i, \sigma^j] = 2i\epsilon_{ijk}\sigma^k, \quad \{\sigma^i, \sigma^j\} = 2\delta_{ij}, \quad (11)$$

where  $i, j, k$  can be 1, 2, or 3 for each possible axis, we see that different pauli matrices do not commute and will actually anticommute. So for instance, the term  $\sigma_1^x \sigma_2^x$  will commute with other terms acting on the same two sites such as  $\sigma_1^y \sigma_2^y$  but will not commute with terms such as  $\sigma_2^z \sigma_3^z$ .

Because our system is only three qubits, we can easily write out the full time evolution operator

$$U(t) = \exp \left[ -i(\sigma_1^x \sigma_2^x + \sigma_1^y \sigma_2^y + \sigma_1^z \sigma_2^z + \sigma_2^x \sigma_3^x + \sigma_2^y \sigma_3^y + \sigma_2^z \sigma_3^z)t \right]. \quad (12)$$

With trotterization over  $N$  time steps this becomes

$$\begin{aligned} U(t) &\approx \left( e^{\frac{-it}{N} \sum_\gamma \sigma_1^\gamma \sigma_2^\gamma} e^{\frac{-it}{N} \sum_\gamma \sigma_2^\gamma \sigma_3^\gamma} \right)^N \\ &\approx \left( \prod_\gamma e^{\frac{-it}{N} \sigma_1^\gamma \sigma_2^\gamma} \prod_\gamma e^{\frac{-it}{N} \sigma_2^\gamma \sigma_3^\gamma} \right)^N, \end{aligned}$$

where the sums/products are over  $\gamma = x, y, z$ . No additional approximation has been made between the first and second lines because the pauli strings acting on the same sites commute. We have now simplified the time evolution operator into unitary gates that we can directly implement on the quantum device. In figure 3 we show an example of a trotterized circuit with  $N = 2$ . Note that this circuit has not been compiled for our specific device, so when the circuit is executed the two qubit gates will need to be decomposed into gates that are in our gate set.

## SIMULATION RESULTS

It is difficult to obtain access to the ibmq-jakarta device, so we will first optimize our circuit using noise models on simulators. Our goal is to simulate the Heisenberg model for a total time  $\tau = \pi$  with the highest fidelity possible. This amount of time  $\tau$  is convenient because our specific initial condition  $|\psi(0)\rangle = |110\rangle$  returns exactly to its initial state  $|\psi(\tau)\rangle = |110\rangle$  after this length of time. Before we can compute the quantum fidelity of the final state on our circuit, we must first determine what the state is. On a real device, we do not have direct access to the exact state and must reconstruct the state based on measurement outcomes. So, we perform measurements in the pauli basis and use built-in Qiskit routines to do quantum tomography [citation].

Let  $\tilde{\rho}$  be the state reconstruction output from our quantum tomography process and let  $\rho = |110\rangle\langle 110|$  be



simulator was  $0.45 \pm .001$ . Given the small uncertainties on both fidelities, we can conclude the real device's performance is substantially worse than the noisy simulator and that this performance difference statistically significant.

## CONCLUSION

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- [1] L. Dupays, arXiv preprint arXiv:2107.01204 (2021).