

# Finding differential trails using an evolutionary algorithm

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Lunch colloquium  
December 7, 2018





# Outline

Background information

Finding trails

Results

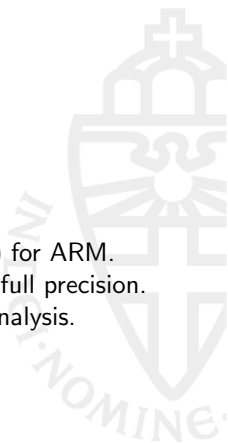
Conclusion





# Research internship

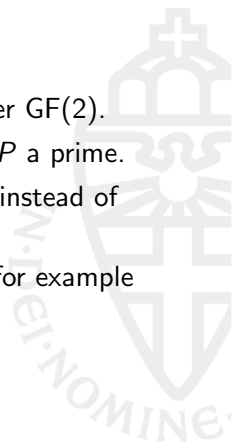
- Part of the TRU/e master.
- 2.5 days a week for one semester.
- Supervised by Joan Daemen.
- Overarching theme: cryptography over  $GF(P)$ 
  - Bitsliced  $+$ ,  $\cdot$  and  $^2$  over  $GF(3)$ ,  $GF(5)$  and  $GF(7)$  for ARM.
  - Multidimensional discrete Fourier transform with full precision.
  - Finding MAC collisions though differential cryptanalysis.
  - Key recovery after finding a collision.





# Cryptography in $\text{GF}(P)$ - What?

- Normally, in symmetric cryptography we work over  $\text{GF}(2)$ .
- Here, we work over a Galois Field of size  $P$  with  $P$  a prime.
- This means we work with digits in  $\{0, \dots, P - 1\}$  instead of bits that are either 0 or 1.
- In  $\text{GF}(P)$  things often behave a little differently, for example addition no longer is the same as subtraction.





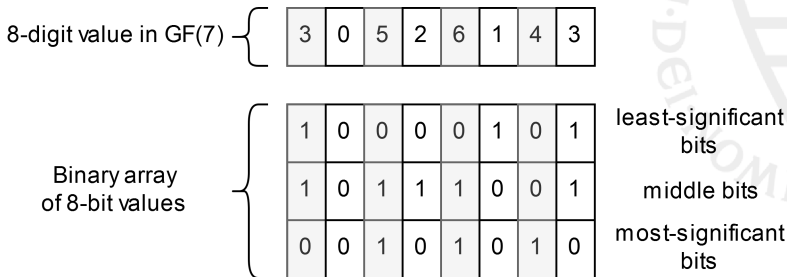
# Cryptography over $GF(P)$ - Why?

- Not only functions over  $GF(2)$  and  $GF(2^n)$  have been studied extensively.
- Interesting differences that to a certain extent provide richer functionality than binary fields.
  - $GF(P)$ : Perfect Nonlinear Functions.
  - $GF(2)$ : Almost Perfect Nonlinear Functions.
- Try to capitalize on the large amount of theoretical results over  $GF(P)$ .
- Generalization might give new insights.
- Applicable to coding theory and signal modulation techniques?



# Cryptography over $GF(P)$ - How?

- Specialized hardware - ternary computers are virtually nonexistent, but again, it might be useful in telecommunications.
- In practice, most data will continue to be represented in bit strings.
- Software emulation -  $\lceil \log_2 P \rceil$  bits required to store a digit in  $GF(P)$ .





# Bitsliced operations for ARM

Bitslicing is a technique where a computation is:

- ① reduced to elementary operations (e.g. OR, AND, XOR);
- ② executed in parallel, with as many simultaneous instances as there are bits in a register.

To work, data needs to be transposed:

Normal				Bitsliced			
$r_0$	$a_2$	$a_1$	$a_0$	$r_0$	$c_0$	$b_0$	$a_0$
$r_1$	$b_2$	$b_1$	$b_0$	$r_1$	$c_1$	$b_1$	$a_1$
$r_2$	$c_2$	$c_1$	$c_0$	$r_2$	$c_2$	$b_2$	$a_2$



# Bitsliced operations for ARM

Normal				Bitsliced			
$r_0$	$a_2$	$a_1$	$a_0$	$r_0$	$c_0$	$b_0$	$a_0$
$r_1$	$b_2$	$b_1$	$b_0$	$r_1$	$c_1$	$b_1$	$a_1$
$r_2$	$c_2$	$c_1$	$c_0$	$r_2$	$c_2$	$b_2$	$a_2$

For values  $a$ ,  $b$  and  $c$ , compute in parallel:  $X_0 + X_0 \cdot X_1$ .

AND      R3, R0, R1      ; R3 := X0 \* X1

EOR      R3, R0, R4      ; R3 := X0 + (X0 \* X1)

Can sometimes be used to dramatically increase throughput at the cost of increased latency.





# Bitsliced operations for ARM

- Addition, multiplication and squaring in  $GF(3)$ ,  $GF(5)$  and  $GF(7)$ .
- Represent the output in terms of a function of the input bits.
- Difficult to do by hand  $\rightarrow$  use logic synthesis tools.
- Implemented in assembly for ARM Cortex M4.

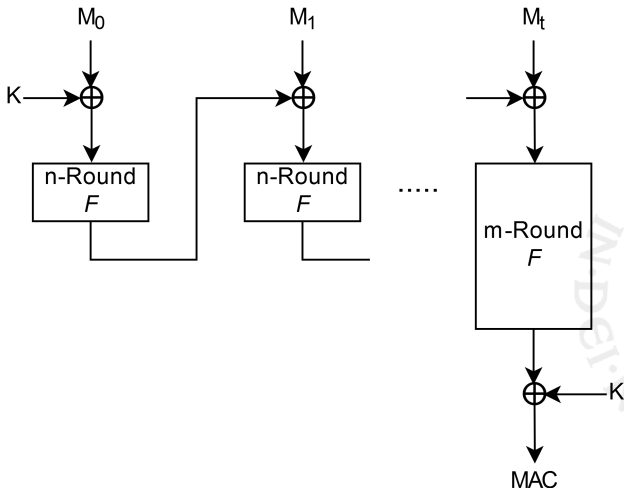


# Multidimensional discrete Fourier transform

- $GF(P)$  version of the Walsh-Hadamard transform.
- Important in linear cryptanalysis.
- Used to compute the correlation of a function with all linear functions at once.
- In the binary case, correlation is between  $-1$  and  $1$ .
- In the non-binary case, it deals with complex numbers.
- Two versions:
  - a straightforward one that uses floating points. Numbers are represented as  $a + bi$ .
  - one that uses polynomials of form  $\sum a_j \omega^j$  with  $\omega = e^{\frac{2\pi i}{P}}$
- The latter has full precision.



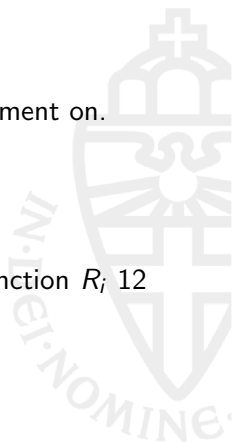
# MACs using (round reduced) transformations





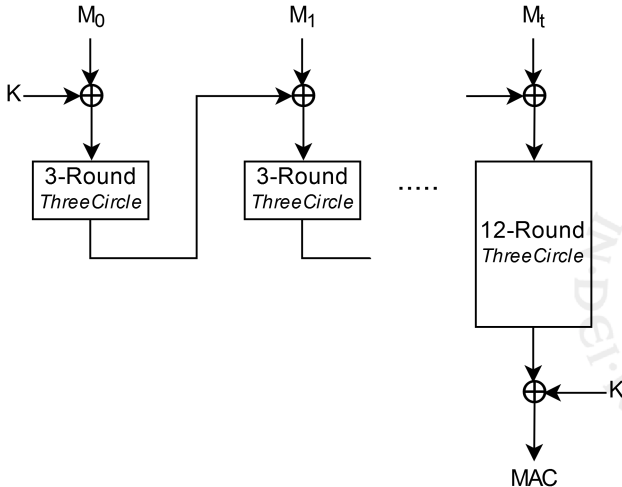
# The cryptographic transformation ThreeCircle

- A permutation designed by Joan for me to experiment on.
- Operates on a state of 160 digits over  $\text{GF}(3)$ .
- State  $a$  arranged in 5 32-digit lanes, we write  $a = (a_0, a_1, a_2, a_3, a_4)$ .
- Classical iterated structure: it iterates a round function  $R_i$  12 times to the state.





# MACs using (round reduced) ThreeCircle





# The round function of ThreeCircle

## Non-linear step $\chi$

- $a_y \leftarrow a_y + a_{y+1}a_{y+1} \lll 1$  for all  $y$  in parallel

## Mixing step $\theta$

- $p \leftarrow a_0 + a_1 + a_2 + a_3 + a_4$
- $e \leftarrow p \lll 12 + p \lll 17$
- $a_y \leftarrow a_y + (e \lll y)$  for all  $y$

## Transposition step $\pi$

- $a_y \leftarrow a_{y+1}$  for all  $y$  in parallel

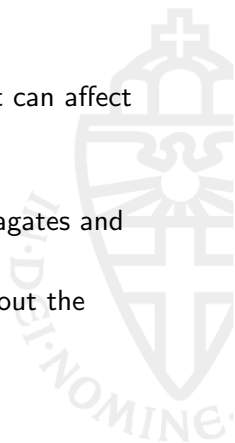
## Transposition step $\rho$

- $a_y \leftarrow a_y \lll r_y$  for all  $y$  with  $r = (0, 2, 6, 11, 19)$



# Differential cryptanalysis

- The study of how differences in information input can affect the resultant difference at the output.
- We want to use it to find colliding MACs.
- Trail: knowledge of how an input difference propagates and the probability of it happening.
- Differential: same as trail, but we do not care about the difference between each round.

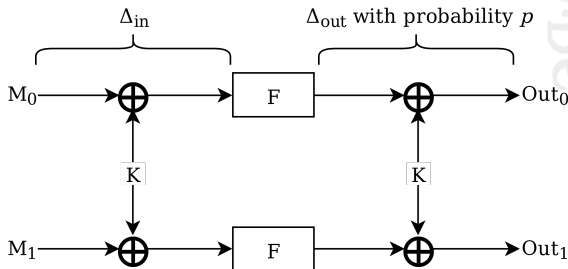




# Differential cryptanalysis - Differentials

- Assume we have the following differential:  $(\Delta_{in}, \Delta_{out})$  with associated differential probability  $DP(\Delta_{in}, \Delta_{out}) = p$ .
- We start with two messages such that  $\Delta_{in} = M_1 - M_0 = (M_1 + K) - (M_0 + K)$
- With probability  $p = DP(\Delta_{in}, \Delta_{out})$ , we have that:  

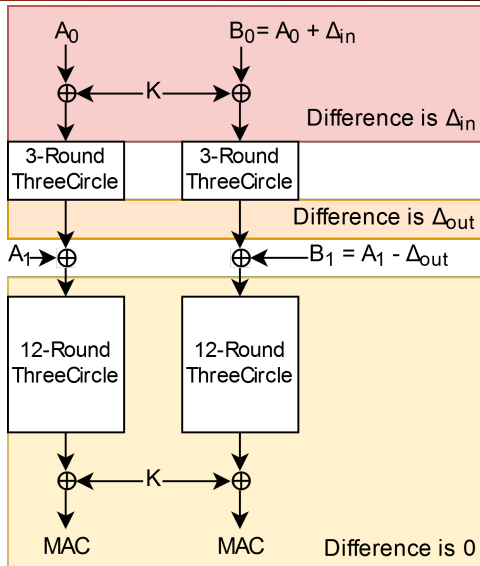
$$\Delta_{out} = F(M_1 + K) - F(M_0 + K) = (F(M_1 + K) + K) - (F(M_0 + K) + K)$$







# Differential cryptanalysis - Finding collisions





# How do differences propagate in ThreeCircle?

## Non-linear layer

- Each non-zero digit in the differential causes 3 equally likely branches.
- Branches lower the differential probability.

## Mixing layer

- Only has effect when there are non-zero digits in differential of the parities.
- Does not cause branching, but causes diffusion (i.e. more non-zero digits in the differential).

## Transposition layer

- Moves the digits around.
- Does not cause branching, but spreads the active digits.



# Finding trails

- Finding best trails is an open problem.
- Finding a trail that spans 1 round often is doable by hand.
- Finding a trail that spans multiple rounds gets *very* complex *very* fast.
- Essentially an optimization problem: we want to find an input/output difference that has an as high as possible differential probability.
- Problem: search space is *huge*.



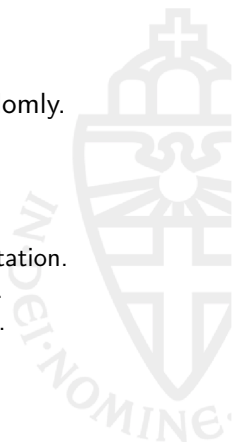
# Evolutionary algorithms

- An evolutionary algorithm (EA) is a generic population-based optimization algorithm.
- An EA uses mechanisms inspired by biological evolution, such as reproduction, mutation, recombination and selection.
- Candidate solutions play the role of individuals in a population, and the fitness function determines the quality of the solutions.
- It often works well on NP-problems.
- Based on underlying assumptions that are not necessarily true, most notably the building block hypothesis.



# Evolutionary algorithms - Implementation

- ① Generate the initial population of individuals randomly.
- ② Evaluate the fitness of each individual.
- ③ Repeat until termination:
  - Select the best-fit individuals for reproduction.
  - Breed new individuals through crossover and mutation.
  - Evaluate the individual fitness of new individuals.
  - Replace least-fit population with new individuals.





# Evolutionary algorithms - fitness function

- Many design decisions when implementing an EA.
- Perhaps the most important one is the fitness function.
- Single figure of merit that summarizes how close a solution is to achieving a set of aims.
- If chosen poorly, the algorithm will converge on a poor solution or not at all.
- We want to minimize the number of branches  $\rightarrow$  minimize the number of active digits going into  $\chi$ .
- Use sum of weights going to  $\chi$ . Problem: noisy.



# Evolutionary algorithms - Tweaks to make things

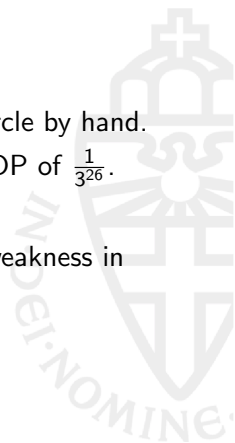
work

- The fitness function is noisy since many trails are possible for a given  $\Delta_{in}$ , therefore I run the fitness function many times and return weight of best trail.
- Empirically, I know  $\Delta_{in}$ 's of good trails tend to have most digits set to 0, so to speed up the process I create the initial population such that most digits are 0.



# Results

- I could not find a good trail for 3 round ThreeCircle by hand.
- We anticipated that the best trail would have a DP of  $\frac{1}{3^{26}}$ .
- I used an EA to find a trail with DP  $\frac{1}{3^{16}}$ .
- The trail I found should not exist and reveals a weakness in (the rotation constants of) ThreeCircle.







# Conclusion

- It might be worth pursuing the use of evolutionary algorithms to find trails for differential cryptanalysis.

