CPSC 313 Spring 2016

Decidability and reductions

- **1.** Show that \leq_m is a transitive relation.
- **2.** Let $T = \{ \langle M \rangle | M \text{ is a Turing machine that accepts } w^R \text{ whenever it accepts } w \}$. Show that T is undecidable.
- **3.** Consider the problem of determining whether a two-tape Turing machine M ever writes a non-blank symbol on its second tape when it is run on some input w. Formulate this problem as a language and show that it is undecidable.
- **4.** Consider the problem of determining whether a two-tape Turing machine *M* ever writes a non-blank symbol on its second tape during the course of its computation on any input string. Formulate this problem as a language and show that it is undecidable.
- **5.** Consider the problem of testing whether a Turing machine M on an input w ever attempts to move its head left when its head is on the left-most tape cell. Formulate this problem as a language and show that it is undecidable.