

CPSC 331

T#5

Worse Case Running Time

Public static Boolean distinctEntires (int[] A){	cost	time
for (int i=1;i<A.lenght;i++) {	c1	n
for (int j = 0; j<i; j++){	c2	(n-1)(n) 2+3+...+n
if (A[j] == A[i]) {	c3	(n-1)(n-1) 1+2+3+...+(n-1)
return false;	c4	1
};		
};		
};		
return true;	c4	1
}		

- In order to ease our analysis of the procedure:
 - Ignore the actual cost of each statement, c_i
These constants give us more detail than we really need
 - We need one more simplification
the order of growth of the running time really interests us . So consider only the leading term of a formula
Since the lower-order terms are relatively insignificant for large values of n .
 - Ignore the leading term's constant coefficient

Growth of Functions

- ❖ For a given function $g(n)$ we have :
 - ❖ $\Theta(g(n)) = \{f(n) : \text{there exist positive constants } c_1, c_2 \text{ and } n_0 \text{ such that } 0 \leq c_1 g(n) \leq f(n) \leq c_2 g(n) \text{ for all } n \geq n_0\}$
 - ❖ Function $f(n)$ belongs to the set $\Theta(g(n))$
- ❖ $O(g(n)) = \{f(n) : \text{there exist positive constants } c \text{ and } n_0 \text{ such that } 0 \leq f(n) \leq c g(n) \text{ for all } n \geq n_0\}$
- ❖ To give an upper bound of a function
- ❖ $\Omega(g(n)) = \{f(n) : \text{there exist positive constants } c \text{ and } n_0 \text{ such that } 0 \leq c g(n) \leq f(n) \text{ for all } n \geq n_0\}$

Relations Between Θ , O , Ω

