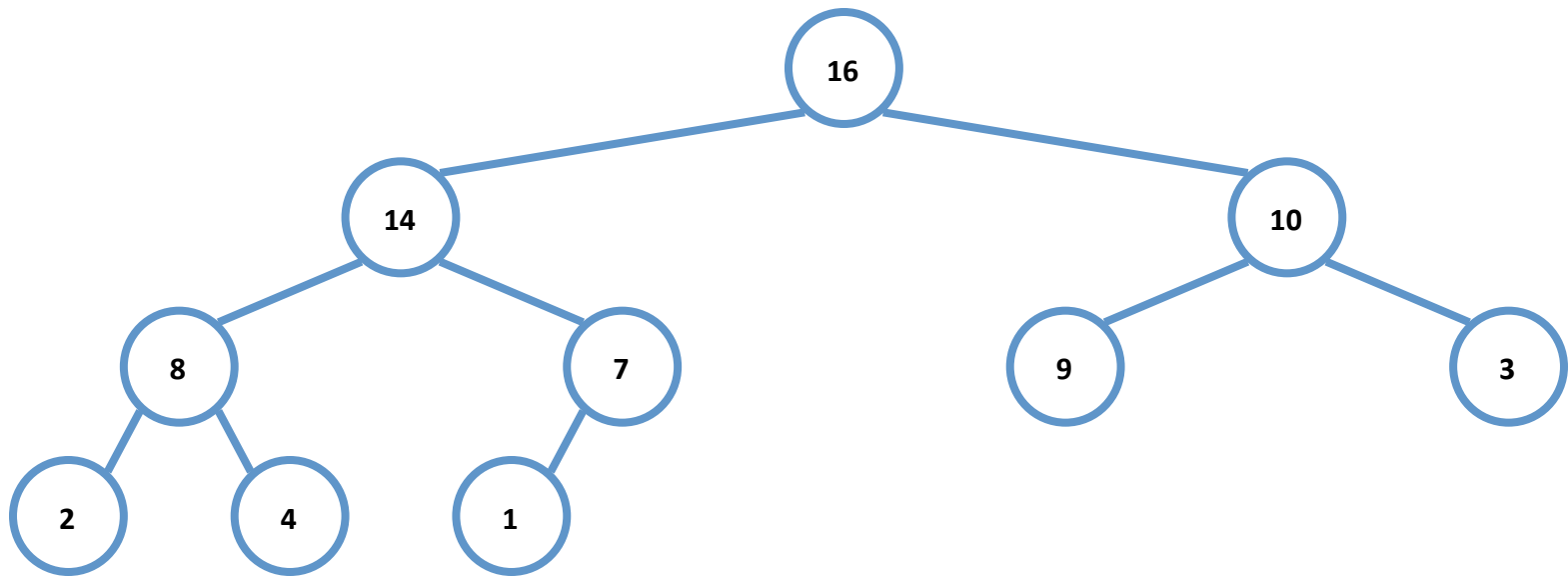


Heap Sort

Heaps

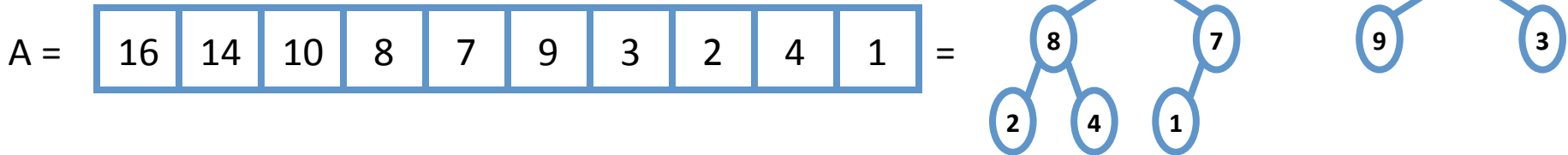
- A *heap* can be seen as a complete binary tree:



- *What makes a binary tree complete?*
- *Is the example above complete?*

Heaps

- In practice, heaps are usually implemented as arrays:



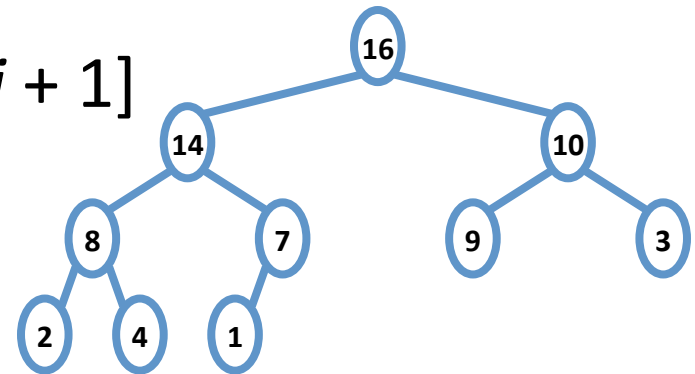
Heaps

- To represent a complete binary tree as an array:
 - The root node is $A[1]$
 - Node i is $A[i]$
 - The parent of node i is $A[i/2]$ (note: integer divide)
 - The left child of node i is $A[2i]$
 - The right child of node i is $A[2i + 1]$

$A =$

16	14	10	8	7	9	3	2	4	1
----	----	----	---	---	---	---	---	---	---

 $=$



Referencing Heap Elements

```
Parent(i) { return  $\lfloor i/2 \rfloor$ ; }
```

```
Left(i) { return  $2*i$ ; }
```

```
right(i) { return  $2*i + 1$ ; }
```

Heap Property

- Heaps also satisfy the *heap property*:

$$A[\textit{Parent}(i)] \geq A[i] \quad \text{for all nodes } i > 1$$

- In other words, the value of a node is at most the value of its parent

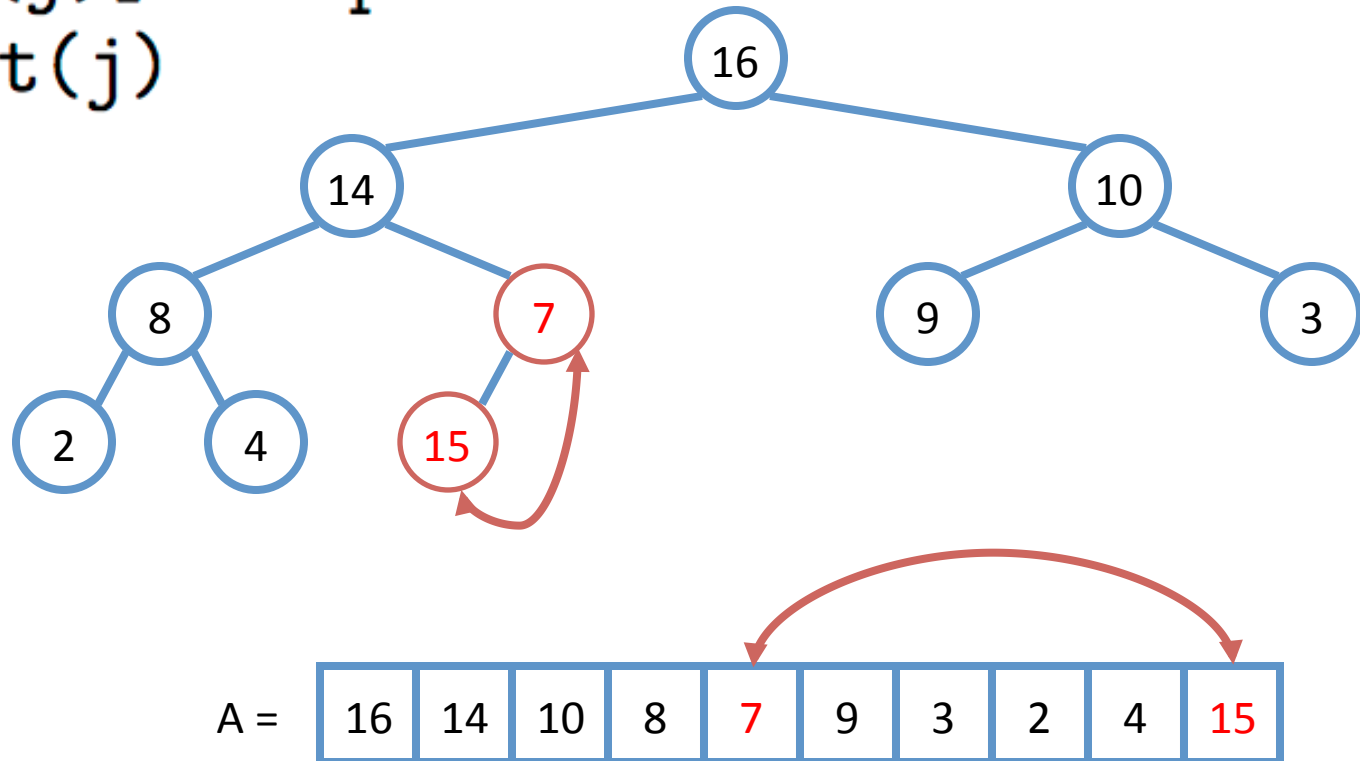
Insert

Pseudocode:

```
void insert(T[] A, T key)
    if heap-size(A) < A.length then
        A[heap-size(A)] = key
        heap-size(A) = heap-size(A) + 1
        The rest of this operation will be described in Step 2
    else
        throw new FullHeapException
    end if
```

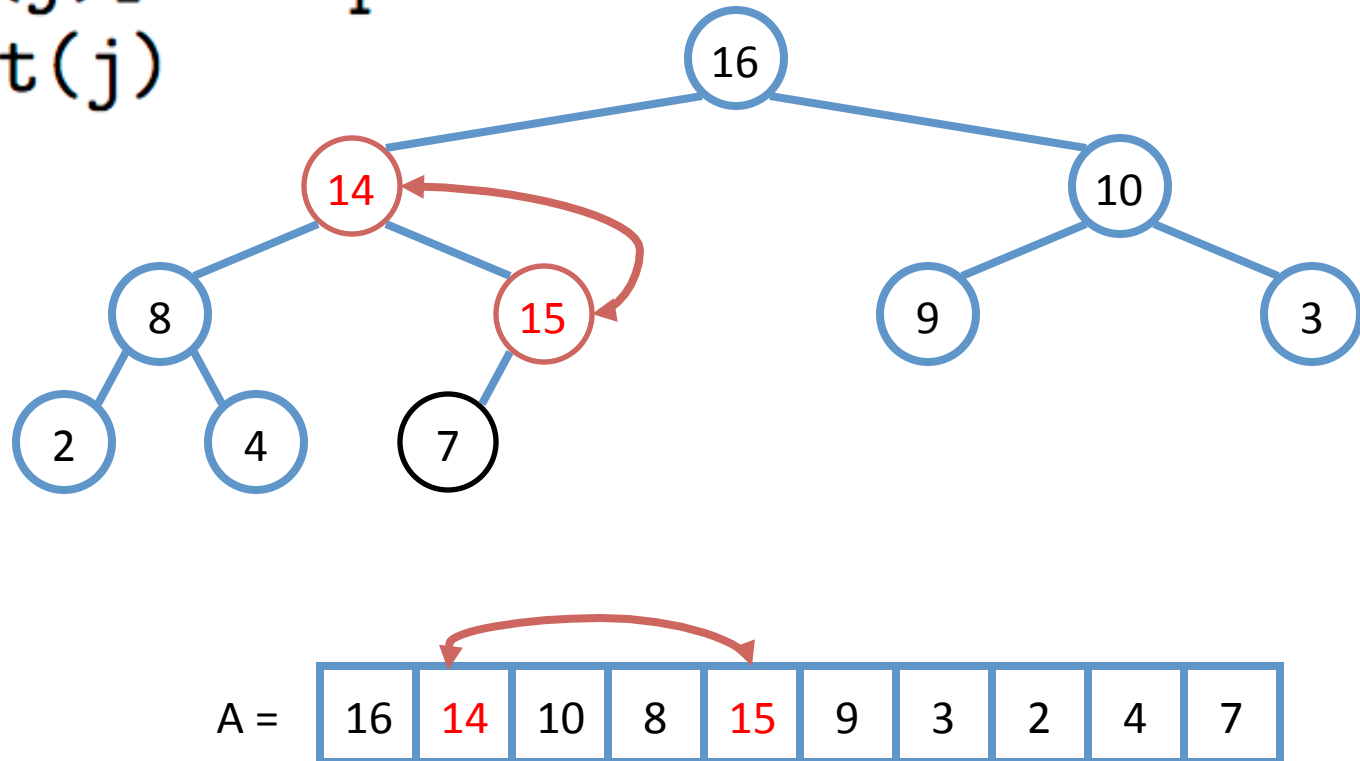
Pseudocode for Step 2:

```
j = heap-size(A) - 1  
while j > 0 and A[j] > A[parent(j)] do  
    tmp = A[j]  
    A[j] = A[parent(j)]  
    A[parent(j)] = tmp  
    j = parent(j)  
end while
```



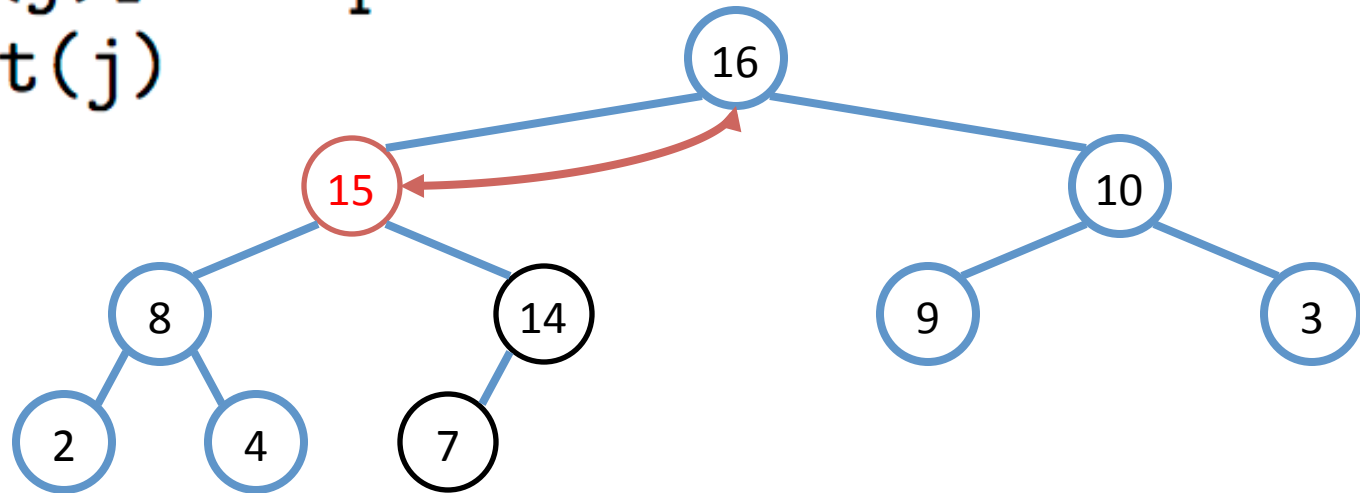
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```



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end while
```

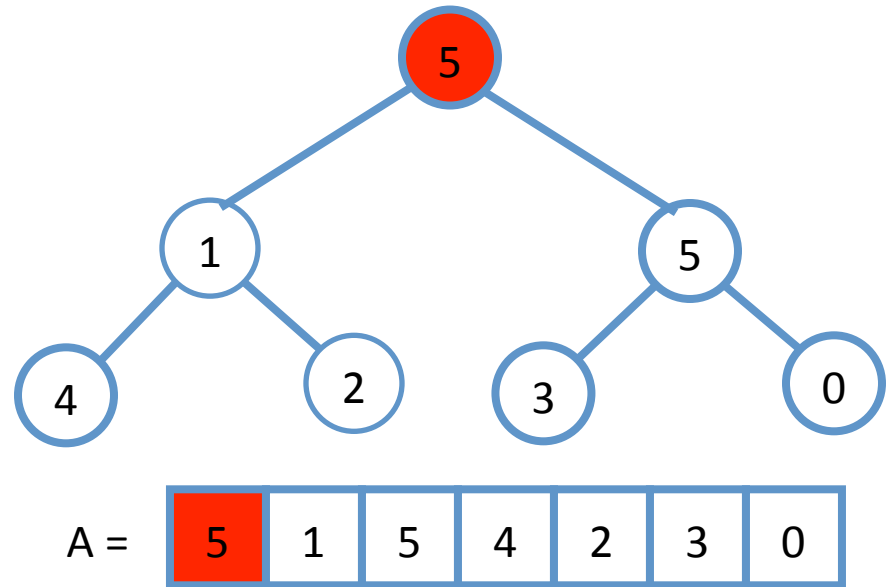


A =

16	15	10	8	14	9	3	2	4	7
----	----	----	---	----	---	---	---	---	---

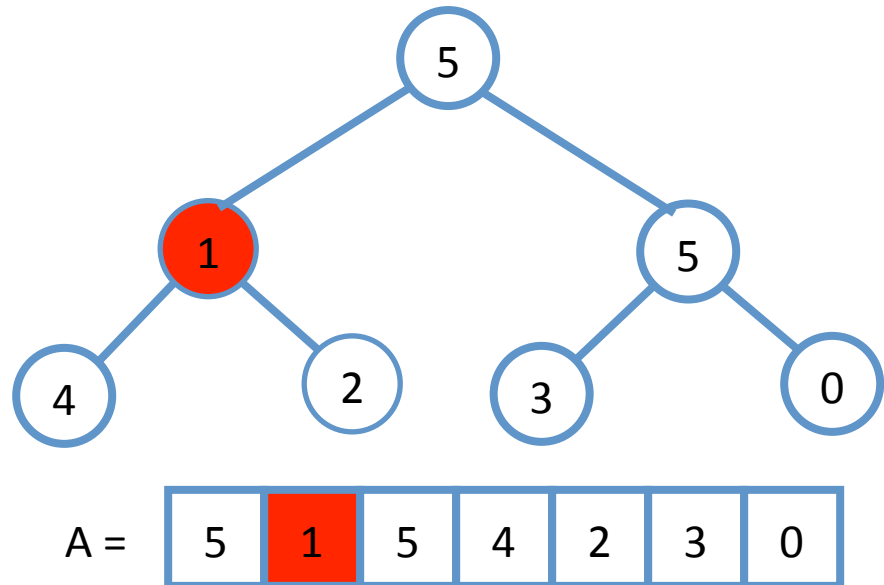
Heap Sort

```
void heapSort(T[] A)
  heap-size(A) = 1
  i = 1
  while i < A.length do
    insert(A, A[i])
    i = i + 1
  end while
  i = A.length - 1
  while i > 0 do
    largest = deleteMax(A)
    A[i] = largest
    i = i - 1
  end while
```



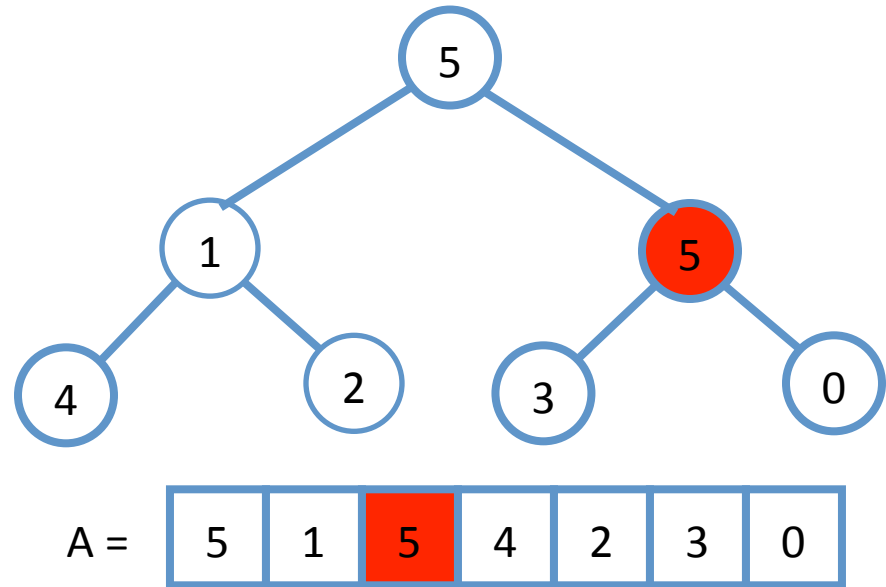
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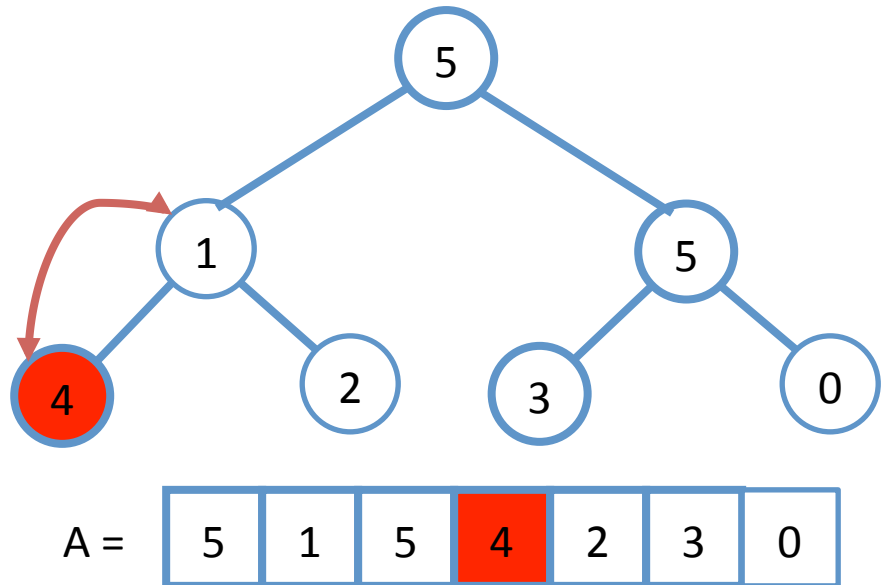
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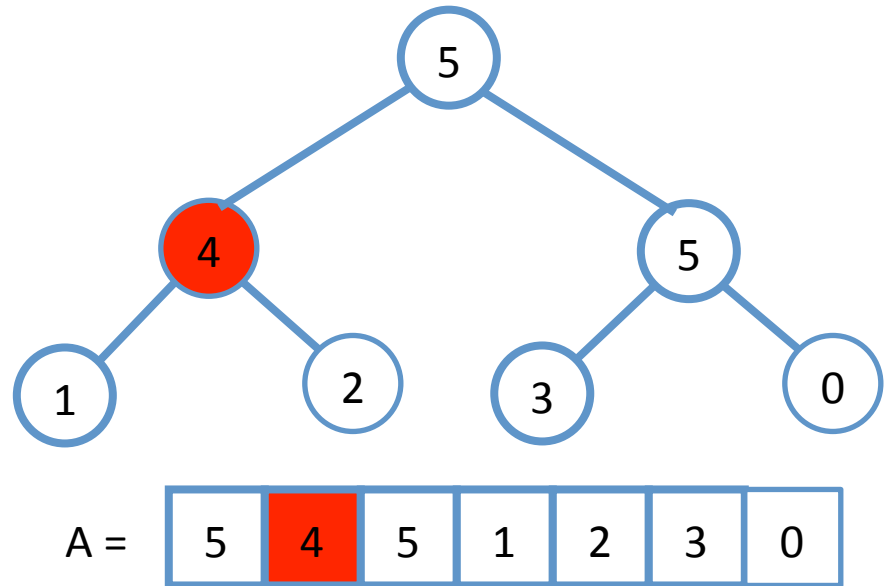
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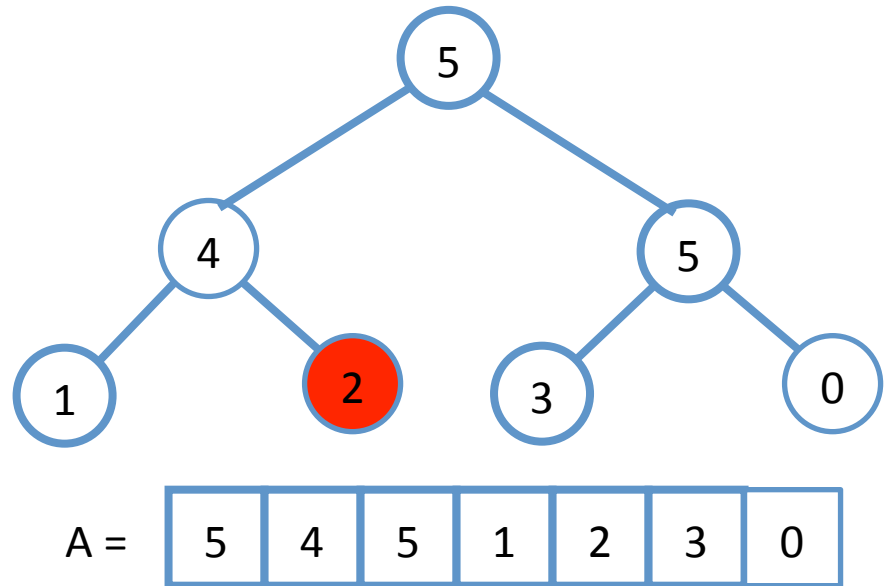
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```



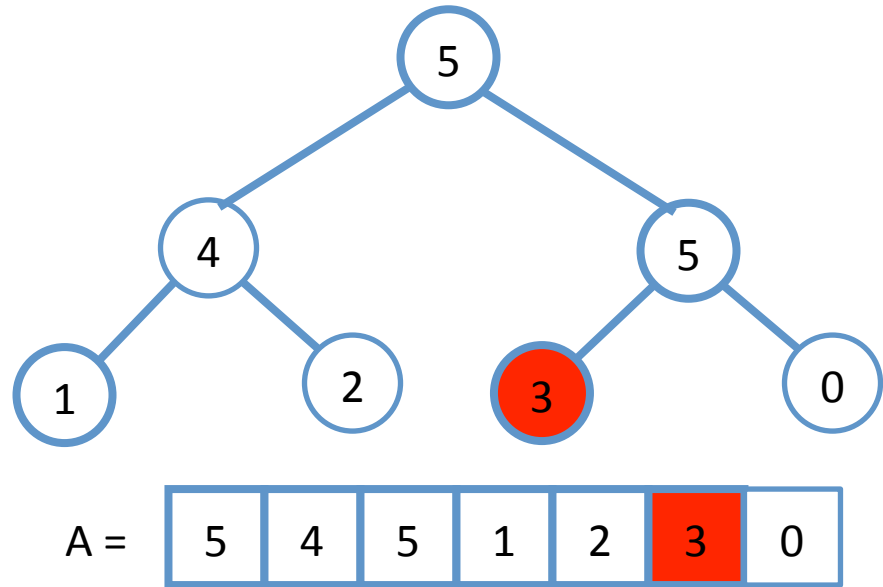
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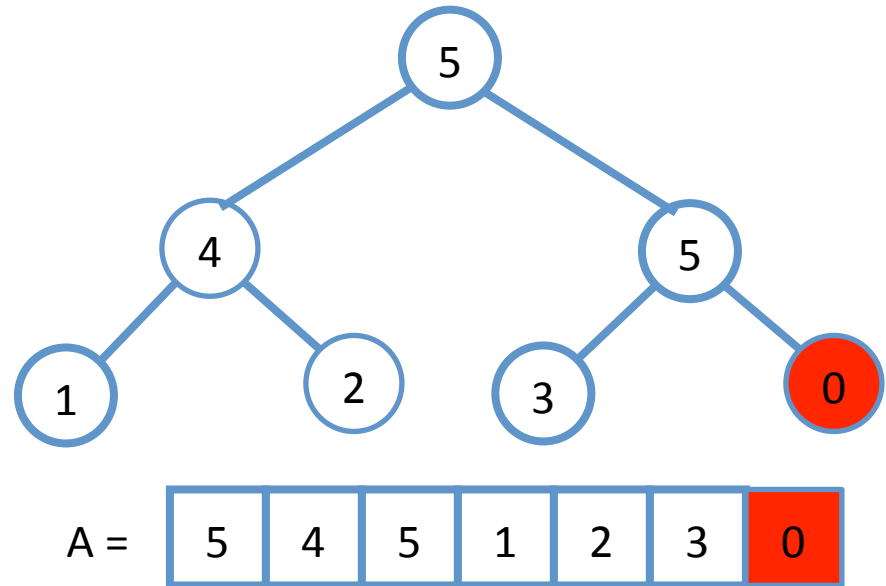
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```



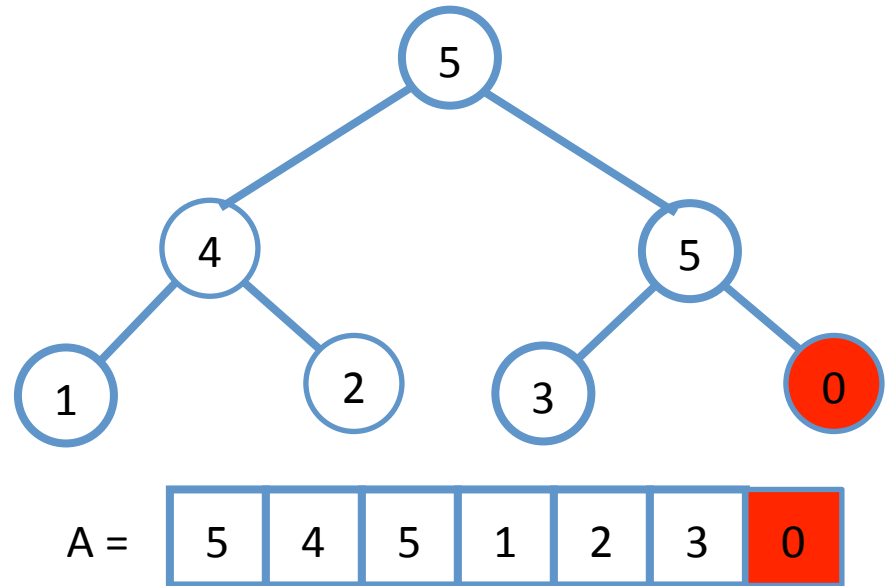
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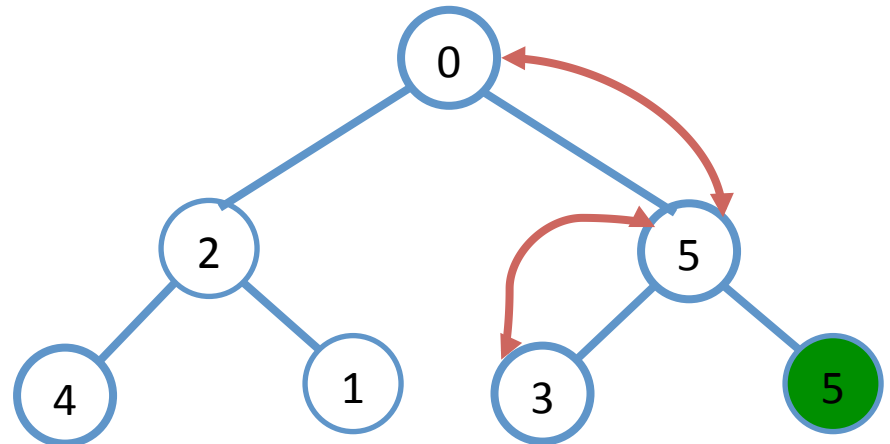
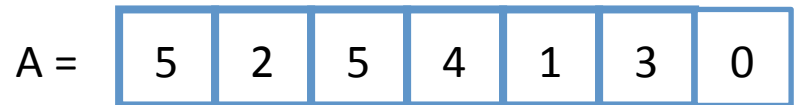
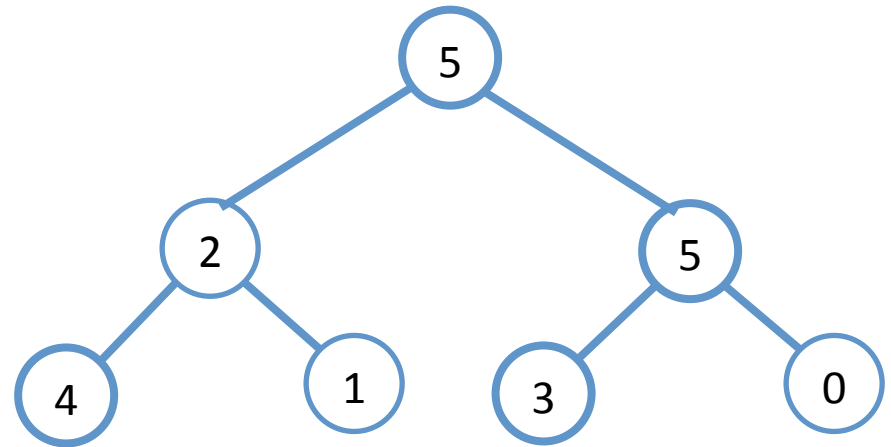
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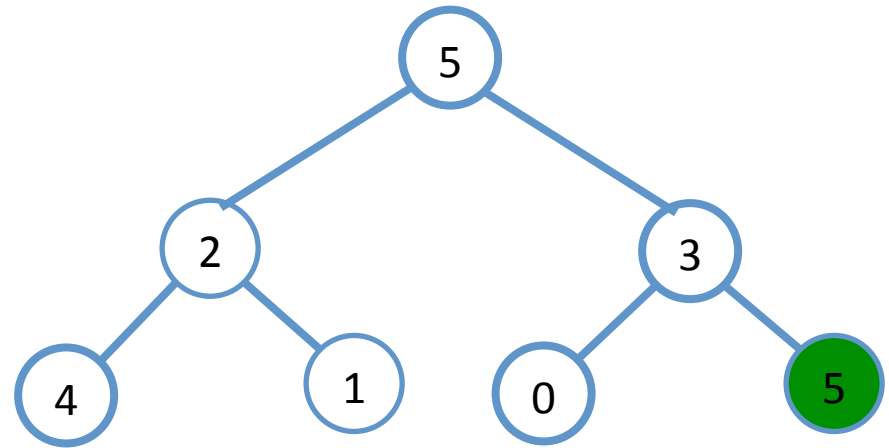
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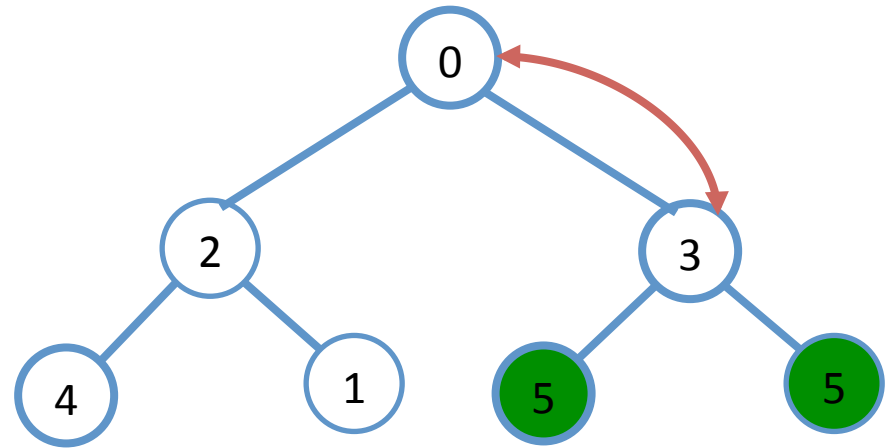
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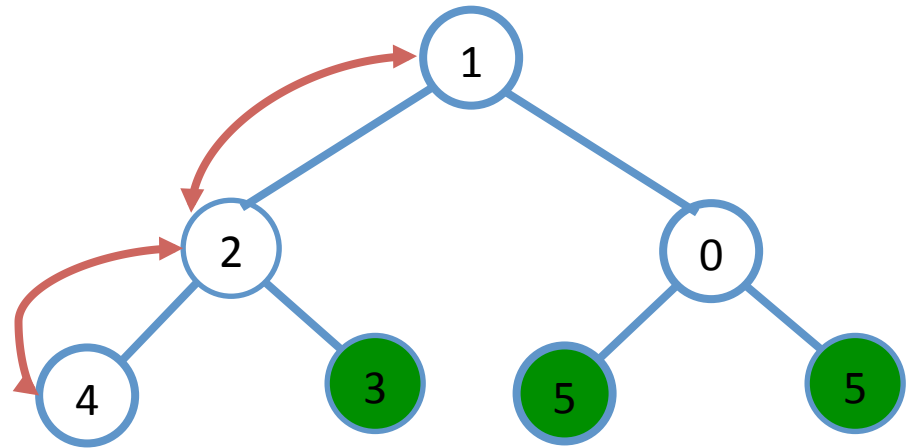
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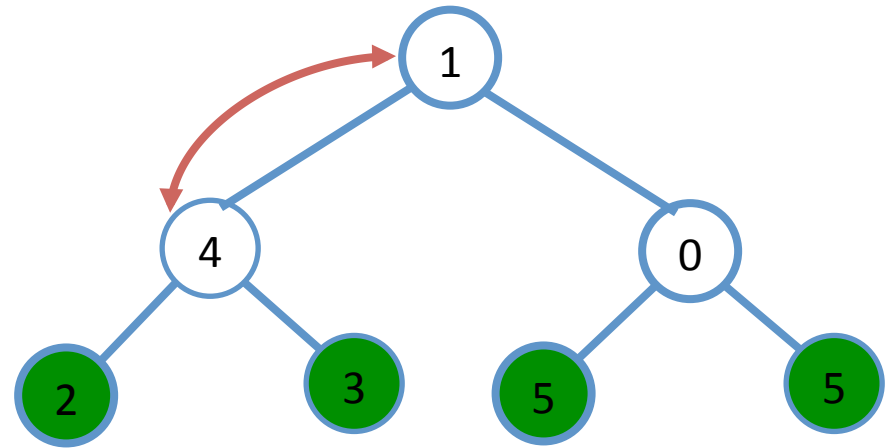
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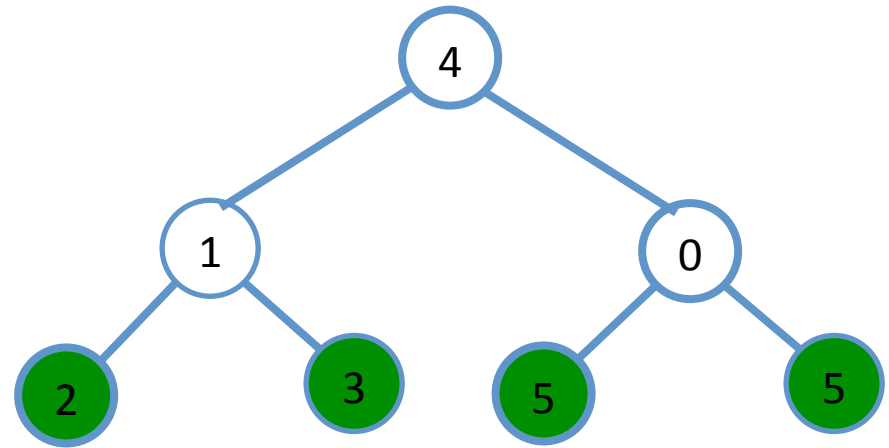
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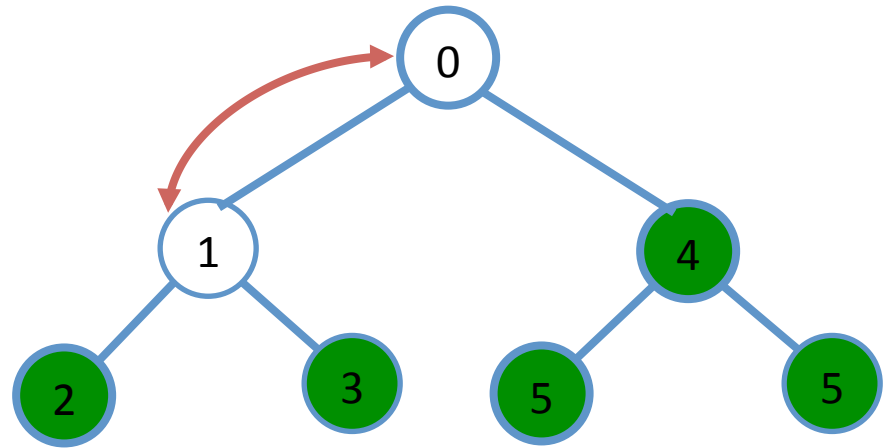
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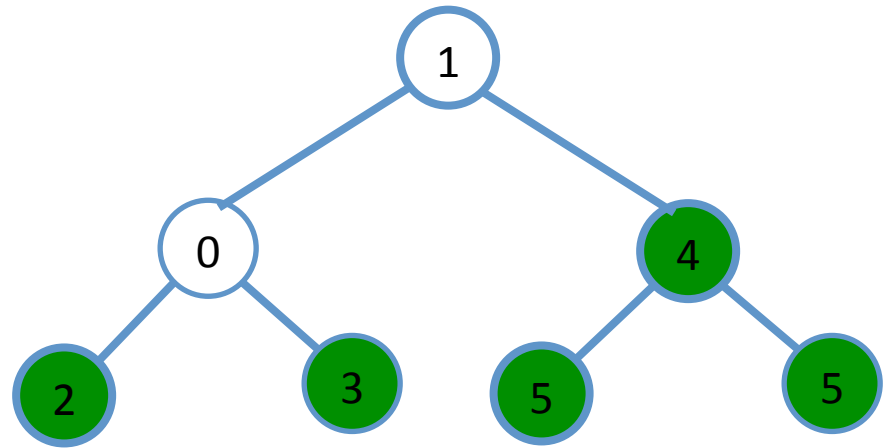
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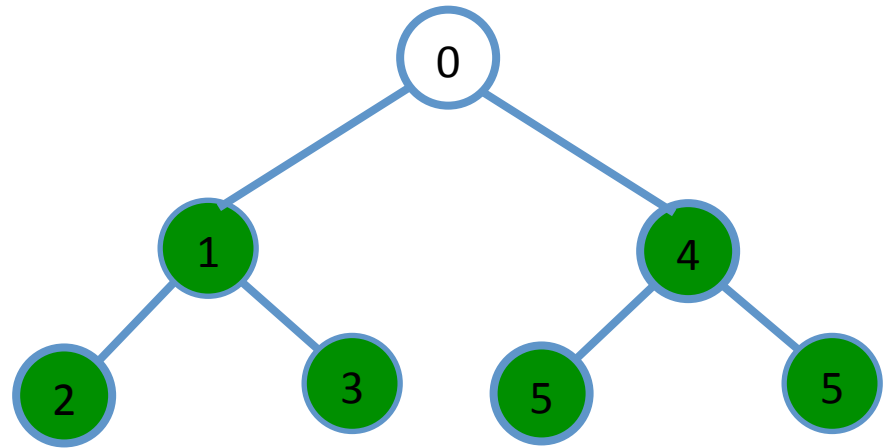
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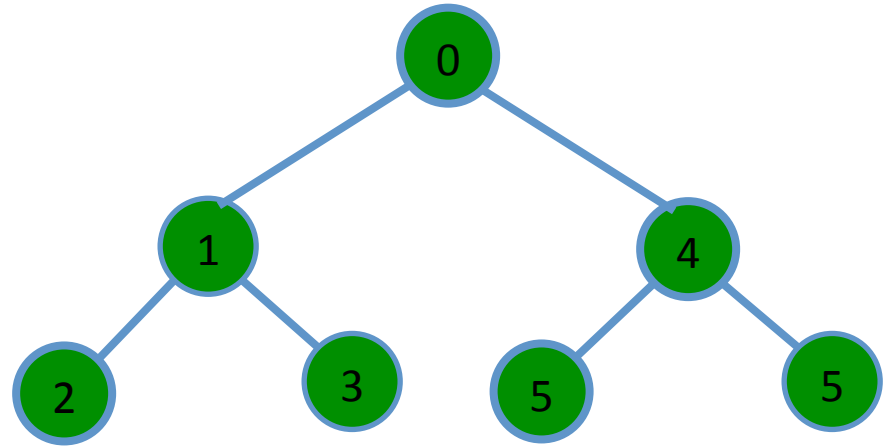
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    end while
```



Build Max Heap Recursively

- Build-Max-Heap(A,i)
 - $N = A.length()$
 - If $Left_child(i) \leq n$ then
 - Build-Max-Heap(A, $Left_child(i)$)
 - If $right_child(i) \leq n$ then
 - Build-Max-Heap(A, $right_child(i)$)
 - Max-Heapify(A, i)

- Max-Heapify(A, i)

set j equal to i the rest of it as follows:

```
while j < heap-size(A) do  
   $\ell = \text{left}(j)$ ;  $r = \text{right}(j)$ ; largest = j  
  if  $\ell < \text{heap-size}(A)$  and  $A[\ell] > A[\text{largest}]$  then  
    largest =  $\ell$   
  end if  
  if  $r < \text{heap-size}(A)$  and  $A[r] > A[\text{largest}]$  then  
    largest = r  
  end if  
  if largest  $\neq j$  then  
    tmp = A[j]; A[j] = A[largest]; A[largest] = tmp;  
    j = largest  
  else  
    j = heap-size(A)  
  end if  
end while
```

Analyzing the time complexity of recursive Build-Max-Heap

- if $m \geq 2$ then $T(m)$ satisfies the following inequality:
 - $T(m) \leq T(m_L) + T(m_R) + c \log m$
- Prove that:
 - $T(m) \in O(m)$.
- Hint :
- It will be helpful to try to show that, for sufficiently large m ,
- $T(m) \leq c_1 m - c_2 \log m$