

MATHEMATICS 271 L01 FALL 2015
ASSIGNMENT 2 SOLUTION

1. Define the sequences a_0, a_1, a_2, \dots and b_0, b_1, b_2, \dots recursively as follows:
 $a_0 = 0$, and for $n > 0$, $a_n = a_{\lfloor \frac{n}{5} \rfloor} + a_{\lfloor \frac{3n}{5} \rfloor} + n$, and
 $b_0 = 2$, $b_1 = 3$ and for $n > 1$, $b_n = 3b_{n-1} - 2b_{n-2}$.
- (a) Find a_1, a_2, a_3, a_4 and a_5 .
 (b) Prove that $a_n \leq 20n$ for all integers $n \geq 0$.
 (c) Find b_2, b_3, b_4 and b_5 .
 (d) Guess a formula for b_n using part (c).
 (e) Prove by induction on n that your guess in part (d) is correct for all integers $n \geq 1$.

Solution:

(a)

$$\begin{aligned} a_1 &= a_0 + a_0 + 1 = 0 + 0 + 1 = 1 \\ a_2 &= a_0 + a_1 + 2 = 0 + 1 + 2 = 3 \\ a_3 &= a_0 + a_1 + 3 = 0 + 1 + 4 = 4 \\ a_4 &= a_0 + a_2 + 4 = 0 + 3 + 4 = 7 \\ a_5 &= a_1 + a_3 + 5 = 1 + 4 + 5 = 10 \end{aligned}$$

(b) We prove $a_n \leq 20n$ for all integers $n \geq 0$ by induction on n (strong form).

Basis: ($n = 0$)

$$a_0 = 0 \leq 0 = 20 \times 0.$$

Inductive step: Let $k \geq 1$ be an integer and suppose that for all integers m where $0 \leq m < k$,

$$a_m \leq 20m \quad (IH)$$

We want to show that $a_k \leq 20k$.

Now, since $k \geq 1$,

$$\begin{aligned} a_k &= a_{\lfloor \frac{k}{5} \rfloor} + a_{\lfloor \frac{3k}{5} \rfloor} + k \\ &\leq 20 \left\lfloor \frac{k}{5} \right\rfloor + 20 \left\lfloor \frac{3k}{5} \right\rfloor + k && \text{by (IH)} \\ &\leq 20 \frac{k}{5} + 20 \frac{3k}{5} + k && \text{because } \lfloor x \rfloor \leq x \text{ for all } x \in \mathbb{R} \\ &= 4k + 12k + k \\ &= 17k \\ &\leq 20k \end{aligned}$$

Thus, $a_n \leq 20n$ for all integers $n \geq 0$.

(c)

$$\begin{aligned} b_2 &= 3b_1 - 2b_0 = 3 \times 3 - 2 \times 2 = 5 \\ b_3 &= 3b_2 - 2b_1 = 3 \times 5 - 2 \times 3 = 9 \\ b_4 &= 3b_3 - 2b_2 = 3 \times 9 - 2 \times 5 = 17 \\ b_5 &= 3b_4 - 2b_3 = 3 \times 17 - 2 \times 9 = 33 \end{aligned}$$

(d) From part (c), we see that:

$$\begin{aligned}
b_0 &= 2 = 2^0 + 1 \\
b_1 &= 2 = 2^1 + 1 \\
b_2 &= 5 = 2^2 + 1 \\
b_3 &= 9 = 2^3 + 1 \\
b_4 &= 17 = 2^4 + 1 \\
b_5 &= 32 = 2^5 + 1
\end{aligned}$$

Thus, we guess that $b_n = 2^n + 1$ for all integers $n \geq 0$.

(e) We prove by induction on n that $b_n = 2^n + 1$ for all integers $n \geq 0$.

Bases: ($n = 0, 1$)

We have shown $b_0 = 2^0 + 1$ and $b_1 = 2^1 + 1$ in part (d).

Inductive step: Let $k \geq 2$ be an integer and suppose that for all integers m where $0 \leq m < k$,

$$b_m = 2^m + 1 \quad (IH)$$

We want to show that $b_k = 2^k + 1$.

Now, since $k \geq 2$,

$$\begin{aligned}
a_k &= 3b_{k-1} - 2b_{k-2} \\
&= 3(2^{k-1} + 1) - 2(2^{k-2} + 1) \quad \text{by (IH)} \\
&= 3 \times 2^{k-1} - 2 \times 2^{k-2} + 1 \\
&= 6 \times 2^{k-2} - 2 \times 2^{k-2} + 1 \\
&= 4 \times 2^{k-2} + 1 \\
&= 2^k + 1
\end{aligned}$$

Thus, $b_n = 2^n + 1$ for all integers $n \geq 0$.

2. Define the Fibonacci sequence f_1, f_2, f_3, \dots recursively as follows:

$$f_1 = f_2 = 1, \text{ and for } n \geq 3, f_n = f_{n-1} + f_{n-2}.$$

(a) Prove that for all integers $n \geq 3$, $\gcd(f_n, f_{n-1}) = \gcd(f_{n-1}, f_{n-2})$. (You may want to use Lemma 4.8.2).

(b) Prove that $\gcd(f_n, f_{n-1}) = 1$ for all integers $n \geq 2$.

(c) Prove that $\sum_{i=1}^n f_i^2 = f_{n+1}f_n$ for all integers $n \geq 1$.

(d) Prove that $f_n < \left(\frac{7}{4}\right)^{n-1}$ for all integers $n \geq 2$.

Solution:

(a) Let $n \geq 2$ be an integer.

Lemma 4.8.2 says that if $a = bq + r$ then $\gcd(a, b) = \gcd(b, r)$. We know

$$f_n = f_{n-1} + f_{n-2}. \quad (1)$$

Put $a = f_n$, $b = f_{n-1}$, $q = 1$ and $r = f_{n-2}$. We see that (1) has the form $a = bq + r$, and so by Lemma 4.8.2 we can conclude that $\gcd(f_n, f_{n-1}) = \gcd(f_{n-1}, f_{n-2})$.

(b) We prove this by (normal) induction on n .

Basis: ($n = 1$)

$$\gcd(f_2, f_1) = \gcd(1, 1) = 1.$$

Inductive step: Let $k \geq 1$ be an integer and suppose that

$$\gcd(f_k, f_{k-1}) = 1. \quad [IH]$$

We want to prove that $\gcd(f_{k+1}, f_k) = 1$.

Now, by part (a) we have

$$\begin{aligned} \gcd(f_{k+1}, f_k) &= \gcd(f_k, f_{k-1}) \\ &= 1 \quad \text{by } [IH] \end{aligned} \quad .$$

Thus, $\gcd(f_n, f_{n-1}) = 1$ for all integers $n \geq 2$.

(c) We prove this by (normal) induction on n .

Basis: ($n = 1$)

$$\sum_{i=1}^1 f_i^2 = 1^2 = 1 = 1 \times 1 = f_2 f_1 = f_{1+1} f_1.$$

Inductive step: Let $k \geq 1$ be an integer and suppose that

$$\sum_{i=1}^k f_i^2 = f_{k+1} f_k \quad (IH)$$

We want to prove that $\sum_{i=1}^{k+1} f_i^2 = f_{k+2} f_{k+1}$.

Now,

$$\begin{aligned} \sum_{i=1}^{k+1} f_i^2 &= \left(\sum_{i=1}^k f_i^2 \right) + f_{k+1}^2 \\ &= f_{k+1} f_k + f_{k+1}^2 \\ &= (f_k + f_{k+1}) f_{k+1} \\ &= f_{k+2} f_{k+1}. \end{aligned}$$

Thus, $\sum_{i=1}^n f_i^2 = f_{n+1} f_n$ for all integers $n \geq 1$.

(d) We prove this by (strong) induction on n .

Basis: ($n = 2, 3$)

$$f_2 = 1 = \frac{4}{4} < \frac{7}{4} = \left(\frac{7}{4}\right)^{2-1}, \text{ and}$$

$$f_3 = f_2 + f_1 = 1 + 1 = 2 = \frac{32}{16} < \frac{49}{16} = \left(\frac{7}{4}\right)^{3-1}$$

Inductive step: Let $k \geq 4$ be an integer and suppose that for all integers m where $2 \leq m < k$, we have

$$f_m < \left(\frac{7}{4}\right)^{m-1} \quad (IH)$$

We want to prove that $f_k < \left(\frac{7}{4}\right)^{k-1}$.

Now, since $k \geq 3$,

$$\begin{aligned}
f_k &= f_{k-1} + f_{k-2} \\
&< \left(\frac{7}{4}\right)^{k-2} + \left(\frac{7}{4}\right)^{k-3} && \text{by (IH), } f_{k-1} < \left(\frac{7}{4}\right)^{k-2} \text{ and } f_{k-2} < \left(\frac{7}{4}\right)^{k-3} \\
&= \left(\frac{7}{4}\right)^{-1} \left(\frac{7}{4}\right)^{k-1} + \left(\frac{7}{4}\right)^{-2} \left(\frac{7}{4}\right)^{k-1} \\
&= \left(\left(\frac{7}{4}\right)^{-1} + \left(\frac{7}{4}\right)^{-2}\right) \left(\frac{7}{4}\right)^{k-1} \\
&= \left(\frac{4}{7} + \frac{16}{49}\right) \left(\frac{7}{4}\right)^{k-1} \\
&= \left(\frac{28}{49} + \frac{16}{49}\right) \left(\frac{7}{4}\right)^{k-1} \\
&= \frac{44}{49} \left(\frac{7}{4}\right)^{k-1} \\
&< \left(\frac{7}{4}\right)^{k-1}
\end{aligned}$$

Thus, $f_n < \left(\frac{7}{4}\right)^{n-1}$ for all integers $n \geq 2$.

3. For any sets A and B , we define the *symmetric difference* $A \triangle B$ by $A \triangle B = (A \cup B) - (A \cap B)$. Note that it is also true that $A \triangle B = (A - B) \cup (B - A)$. Let \mathcal{S} be the statement: “For all sets A, B and C , if $A \subseteq B \cup C$ and $B \subseteq C \cup A$ then $A \triangle B = C$.” and let \mathcal{T} be the statement: “For all sets A, B and C , if $A \triangle B = A \triangle C$ then $B \subseteq C$.”

(a) Is \mathcal{S} true? Prove your answer.

Solution: \mathcal{S} is not true. For example, when $A = B = \emptyset$, and $C = \{1\}$ we have $A = \emptyset \subseteq B \cup C$ and $B = \emptyset \subseteq C \cup A$, but $A \triangle B = \emptyset \neq \{1\} = C$.

(b) Is \mathcal{T} true? Prove your answer.

Solution: \mathcal{T} true. Let A, B and C be sets and suppose that $A \triangle B = A \triangle C$. We prove that $B \subseteq C$. Let $x \in B$. We consider two cases $x \in A$ and $x \notin A$.

Case 1: $x \in A$. Since $x \in A$ and $x \in B$, $x \in A \cap B$ and hence $x \notin A \triangle B$. Since $A \triangle B = A \triangle C$ and $x \notin A \triangle B$, we get $x \notin A \triangle C = (A \cup C) - (A \cap C)$ and therefore $x \notin A \cup C$ or $x \in A \cap C$. However, since $x \in A$, $x \in A \cup C$ and it follows that $x \in A \cap C$ and so $x \in C$.

Case 1: $x \notin A$. Since $x \notin A$ and $x \in B$, $x \in B - A$ and hence $x \in A \triangle B$. Since $A \triangle B = A \triangle C$ and $x \in A \triangle B$, we get $x \in A \triangle C = (A \cup C) - (A \cap C)$ and therefore $x \in A \cup C$, which implies that $x \in A$ or $x \in C$. However, since $x \notin A$, we see that $x \in C$.

(c) Write the converse of \mathcal{S} . Is the converse of \mathcal{S} true? Prove your answer.

Solution: The converse of \mathcal{S} is “For all sets A, B and C , if $A \triangle B = C$ then $A \subseteq B \cup C$ and $B \subseteq C \cup A$.”

The converse of \mathcal{S} is true. Let A, B and C be sets and suppose that $A \triangle B = C$.

First, we prove that $A \subseteq B \cup C$ by contradiction. Suppose that $A \not\subseteq B \cup C$. Then there exists an element $x \in A$ so that $x \notin B \cup C$, that is, $x \notin B$ and $x \notin C$. Since $x \in A$ and $x \notin B$, $x \in A - B \subseteq A \triangle B$. Since we have $x \in A \triangle B$ and $x \notin C$, $A \triangle B \neq C$ which contradicts the assumption that $A \triangle B = C$. Thus, $A \subseteq B \cup C$.

Next, we prove that $B \subseteq C \cup A$ by contradiction. Suppose that $B \not\subseteq C \cup A$. Then there exists an element $x \in B$ so that $x \notin C \cup A$, that is, $x \notin C$ and $x \notin A$. Since $x \in B$

and $x \notin A$, $x \in B - A \subseteq A \Delta B$. Since we have $x \in A \Delta B$ and $x \notin C$, $A \Delta B \neq C$ which contradicts the assumption that $A \Delta B = C$. Thus, $A \subseteq B \cup C$.