Bit-Blasting for SMT-NIA with **AProVE**⋆

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In automated termination provers like our tool AProVE [4], often the need arises
to solve Boolean combinations of constraints in non-linear (integer) arithmetic
to perform a proof step for a successful termination proof. Examples for promi-
nent termination techniques where this is the case are well-founded orders based
on polynomial [3] or matrix interpretations [2]. In order to facilitate this task,
AProVE features a dedicated SMT-solver for SMT-NIA. The approach we are
using at SMT-COMP for the QF_NIA category is based on a reduction to a
SAT problem (see [3] for details) and a subsequent invocation of the SAT solver
MiniSAT[1]. This kind of approach is commonly known as *bit-blasting*.

Since satisfiability of SMT-NIA is undecidable in general, we search for solu-
tions from a finite domain for the variables in the SMT-NIA instances. In case
a satisfying assignment is found by the SAT solver, we return the corresponding
integer solution. If the SAT solver detects unsatisfiability of the generated SAT
instance, we know that there is no solution for the SMT-NIA instance with the
used domain, and we restart the search with an extended domain. In this way,
we obtain a semi-decision procedure for SMT-NIA. We additionally use infor-
mation from global constraints (i.e., atomic top-level assertions like $x \geqslant 42$) of
the SMT-NIA instance to bound the search space for certain variables.

References

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⋆ System description for *SMT-COMP 2014*.