2. Exercise

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1 Exercise 2.1

2.1.1 **Definition 1** (relational completeness) A query language is relationally complete if it is at least as expressive as relational algebra.

Methods with Example table $A(x_1,y_1)$ and $B(y_1)$

method	RA (Example)	keyword	SQL(Example)
selection	$\sigma_{x_1>0}(A)$	WHERE	SELECT * FROM A WHERE A. $x_1 > 0$
projection	$\pi_{x_1}(A)$	SELECT	SELECT x_1 FROM A
cartesian product	$A \times B$	FROM, CROSS	SELECT * FROM A, B
union	$\sigma_{x_1>0}(A)\cup$	UNION	SELECT * FROM A WHERE $x_1 > 0$ UNION
	$\sigma_{x_1<0}(A)$		SELECT * FROM A WHERE $x_1 < 0$
difference	A - A	EXCEPT, NOT IN	SELECT * FROM B WHERE y_1 NOT IN A

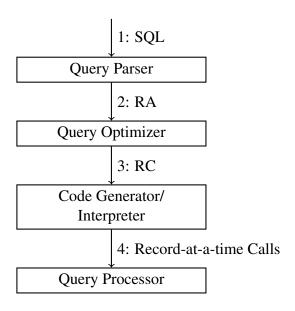
2.1.2 **1. Beispiel**

SELECT SUM(salary) FROM works;

2. Beispiel

SELECT COUNT(city) FROM located;

2.1.3 Solution Graph:



2.1.4 Overview:

operator	operation type	
selection	unary operator	
projection	unary operator	
cartesian product	set & binary operator	
union	set & binary operator	
difference	set & binary operator	
rename	set operator	

2 Exercise 2.2

2.2.1 relational algebra:

 $\pi_{pname}(\rho_p(works) \bowtie_{p.cname=b.cname \land p.pname=b.pname1 \land p.salary > b.salary})$ $\rho_b(works \bowtie_{works.pname=boss.mname} \rho_{pname1 \leftarrow pname}(boss)))$

tuple relational calculus:

 $\{ < w.pname > | \ w \in works \land \exists m \in boss(w.pname = m.pname \land \exists w' \in works(w'.pname = m.mname \land w.cname = w'.cname \land w.salary > w'.salary)) \}$

domain relational calculus:

 $\{pname \mid \exists cname, salary1, salary2, mname \land works(pname, cname, salary1) \land works(mname, cname, boss(pname, cname) \land salary1 > salary2\}$

2.2.2 relational algebra:

 $\pi_{pname}(\rho_{c_1}(works)\bowtie_{c_1.pname=c_2.pname \land c_1.cname!=c_2.cname} \rho_{p_2}(works))$

tuple relational calculus:

 $\{< w_1.pname > | w_1 \in works \land \exists w_2 \in works(w_1.pname = w_2.pname \land w_1.cname \neq w_2.cname)\}$

domain relational calculus:

 $\{pname \mid \exists cname_1, cname_2, salary_1, salary_2 \land works(pname, cname_1, salary_1) \land works(pname, cname_2, salary_2) \land cname_1 \neq cname_2\}$

2.2.3 relational algebra:

 $\pi_{pname}(works) - \pi_{pname}(\rho_{p_1}(works) \bowtie_{p_1.salary < p_2.salary} \rho_{p_2}(works))0$

tuple relational calculus:

 $\{\langle w_1.pname \rangle | w_1 \in works \land \neg(\exists w_2 \in works(w_2.salary))\}$

domain relational calculus:

 $\{pname \mid \exists cname, salary_1 \land works(pname, cname, salary_1) \land \neg(\exists pname_2, salary_2 \land works(pname_2, cname, salary_2) \land salary_2 > salary_1)\}$

2.2.4 relational algebra:

 $\pi_{c_1.cname}(\rho_{c_1}(located) \bowtie_{c_1.citu\neq c_2.citu} \rho_{c_2}(\sigma_{cname='IBM'}(located)))$

tuple relational calculus:

 $\{\langle c_1.cname \rangle | c_1 \in located \land \forall c_2 \in located(c_2.cname \neq' IBM' \lor c_1.city \neq c_2.city)\}$

domain relational calculus:

 $\{cname \mid \exists city \land located(cname, city) \land \neg located('IBM', city)\}$

3 Exercise 2.3

- 2.3.1 with N=15000 and B=8: $\lceil \frac{N}{B} \rceil \Rightarrow \lceil \frac{15000}{8} \rceil = 1875$
- 2.3.2 Number of passes given by: $1 + \lceil \log_{B-1} \lceil \frac{N}{B} \rceil \rceil$ with N = 15000 and $B = 8 \Rightarrow 1 + \lceil \log_{8-1} \lceil \frac{15000}{8} \rceil \rceil = 1 + \lceil \log_7 1875 \rceil = 5$
- $2.3.3 \ \ 2 = 1 + \lceil \log_{B-1} \lceil \frac{15000}{B} \rceil \rceil \Leftrightarrow 1 = \lceil \log_{B-1} \lceil \frac{15000}{B} \rceil \rceil \Leftrightarrow \frac{1}{2} (1 + \sqrt{600001}) \leq B \Rightarrow 123 \leq B$ So we need at-least 123 Buffer to be able to sort the file in two passes.
- 2.3.4 Since we only use 2+1 Pages in a Two-Way-Sort: B=3 and N=15000 $\Rightarrow 1+\lceil \log_{3-1}\lceil \frac{15000}{2}\rceil \rceil=14$ passes and $\lceil \frac{15000}{2}\rceil=7500$ runs