

2. Exercise

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1 Exercise 2.1

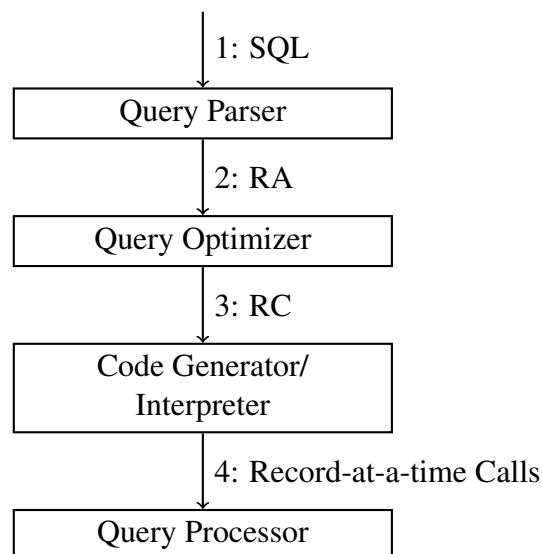
2.1.1 **Definition 1 (relational completeness)** *A query language is relationally complete if it is at least as expressive as relational algebra.*

Methods with Example table $A(\underline{x_1}, x_2, x_3)$ and $B(\underline{y_1}, y_2, y_3)$

method	relational algebra	SQL
selection	$\sigma_{x_1 > 0}(A)$	SELECT * FROM A WHERE A.x ₁ > 0
projection	$\pi_{x_1}(A)$	SELECT x ₁ FROM A
cartesian product	$A \times B$	SELECT * FROM A, B
union	$\sigma_{x_1 > 0}(A) \cup \sigma_{x_1 < 0}(A)$	SELECT * FROM A WHERE A.x ₁ > 0 UNION SELECT * FROM A WHERE A.x ₁ < 0
difference	$A - A$	A MINUS A

2.1.2 TODO

2.1.3 Solution Graph:



2.1.4 Overview:

operator	operation type
selection	unary operator
projection	unary operator
cartesian product	set & binary operator
union	set & binary operator
difference	set & binary operator

2 Exercise 2.2

2.2.1 relational algebra:

$$\pi_{pname}(\rho_p(works) \bowtie_{p.cname=b.cname \wedge p.salary > b.salary} \rho_b(works \bowtie_{works.pname=boss.mname} \rho_{pname1 \leftarrow pname}(boss)))$$

tuple relational calculus:

$$\{ \langle w.pname \rangle \mid w \in works \wedge \exists m \in boss (w.pname = m.pname \wedge \exists w' \in works (w'.pname = m.pname \wedge w.cname = w'.cname \wedge w.salary > w'.salary)) \}$$

domain relational calculus:

$$\{pname \mid \exists cname, salary1, salary2, mname \wedge works(pname, cname, salary1) \wedge works(mname, cname, salary2) \wedge boss(pname, cname) \wedge salary1 > salary2\}$$

2.2.2 relational algebra:

$$\pi_{pname}(\rho_{c_1}(works) \bowtie_{c_1.pname=c_2.pname \wedge c_1.cname \neq c_2.cname} \rho_{p_2}(works))$$

tuple relational calculus:

$$\{ \langle w_1.pname \rangle \mid w_1 \in works \wedge \exists w_2 \in works (w_1.pname = w_2.pname \wedge w_1.cname \neq w_2.cname) \}$$

domain relational calculus:

$$\{pname \mid \exists cname_1, cname_2, salary_1, salary_2 \wedge works(pname, cname_1, salary_1) \wedge works(pname, cname_2, salary_2) \wedge cname_1 \neq cname_2\}$$

2.2.3 relational algebra:

$$\pi_{pname}(works) - \pi_{pname}(\rho_{p_1}(works) \bowtie_{p_1.salary < p_2.salary} \rho_{p_2}(works))$$

tuple relational calculus:

$$\{ \langle w_1.pname \rangle \mid w_1 \in works \wedge \neg (\exists w_2 \in works (w_1.cname = w_2.cname \wedge w_2.salary > w_1.salary)) \}$$

domain relational calculus:

$$\{pname \mid \exists cname, salary_1 \wedge works(pname, cname, salary_1) \wedge \neg (\exists pname_2, salary_2 \wedge works(pname_2, cname, salary_2) \wedge salary_2 > salary_1) \}$$

2.2.4 relational algebra:

$$\pi_{c_1.cname}(\rho_{c_1}(located) \bowtie_{c_1.city \neq c_2.city} \rho_{c_2}(\sigma_{cname='IBM'}(located)))$$

tuple relational calculus:

$$\{ \langle c_1.cname \rangle \mid c_1 \in located \wedge \forall c_2 \in located (c_2.cname \neq 'IBM' \vee c_1.city \neq c_2.city) \}$$

domain relational calculus:

$$\{cname \mid \exists city \wedge located(cname, city) \wedge \neg located('IBM', city) \}$$

3 Exercise 2.3

2.3.1 with $N = 15000$ and $B = 8$: $\lceil \frac{N}{B} \rceil \Rightarrow \lceil \frac{15000}{8} \rceil = 1875$

2.3.2 Number of passes given by: $1 + \lceil \log_{B-1} \lceil \frac{N}{B} \rceil \rceil$
with $N = 15000$ and $B = 8 \Rightarrow 1 + \lceil \log_{8-1} \lceil \frac{15000}{8} \rceil \rceil = 1 + \lceil \log_7 1875 \rceil = 5$

2.3.3 $2 = 1 + \lceil \log_{B-1} \lceil \frac{15000}{B} \rceil \rceil \Leftrightarrow 1 = \lceil \log_{B-1} \lceil \frac{15000}{B} \rceil \rceil \Leftrightarrow \frac{1}{2}(1 + \sqrt{600001}) \leq B \Rightarrow 123 \leq B$
So we need at-least 123 Buffer to be able to sort the file in two passes.

2.3.4 TODO