

CODEFORCES #162

DIV1-E

RODESIDETREES

WRITER: SNUKE

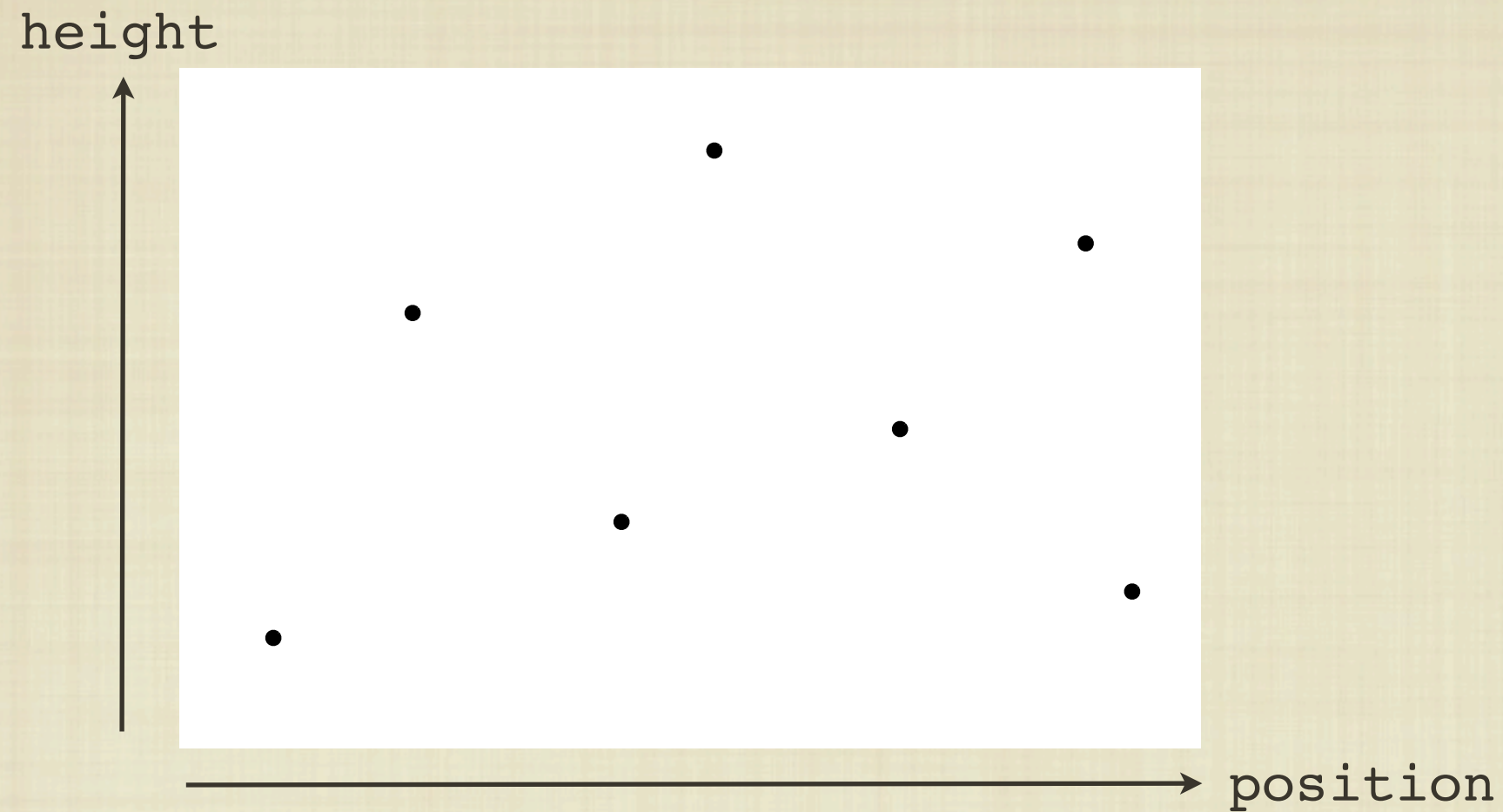
PROBLEM SUMMARY

- You should process queries:
Type1: Plant a tree with height h . ($h \leq 10$)
Type2: Cut the x -th tree from west. ($x \leq 10$)
- After each query, you should calculate the length of LIS of height of tree (from west to east).
- Before each query, trees grow 1 meter.
- N: The number of trees $\leq 10^5$
M: The number of queries $\leq 2 \cdot 10^5$

SOLUTION

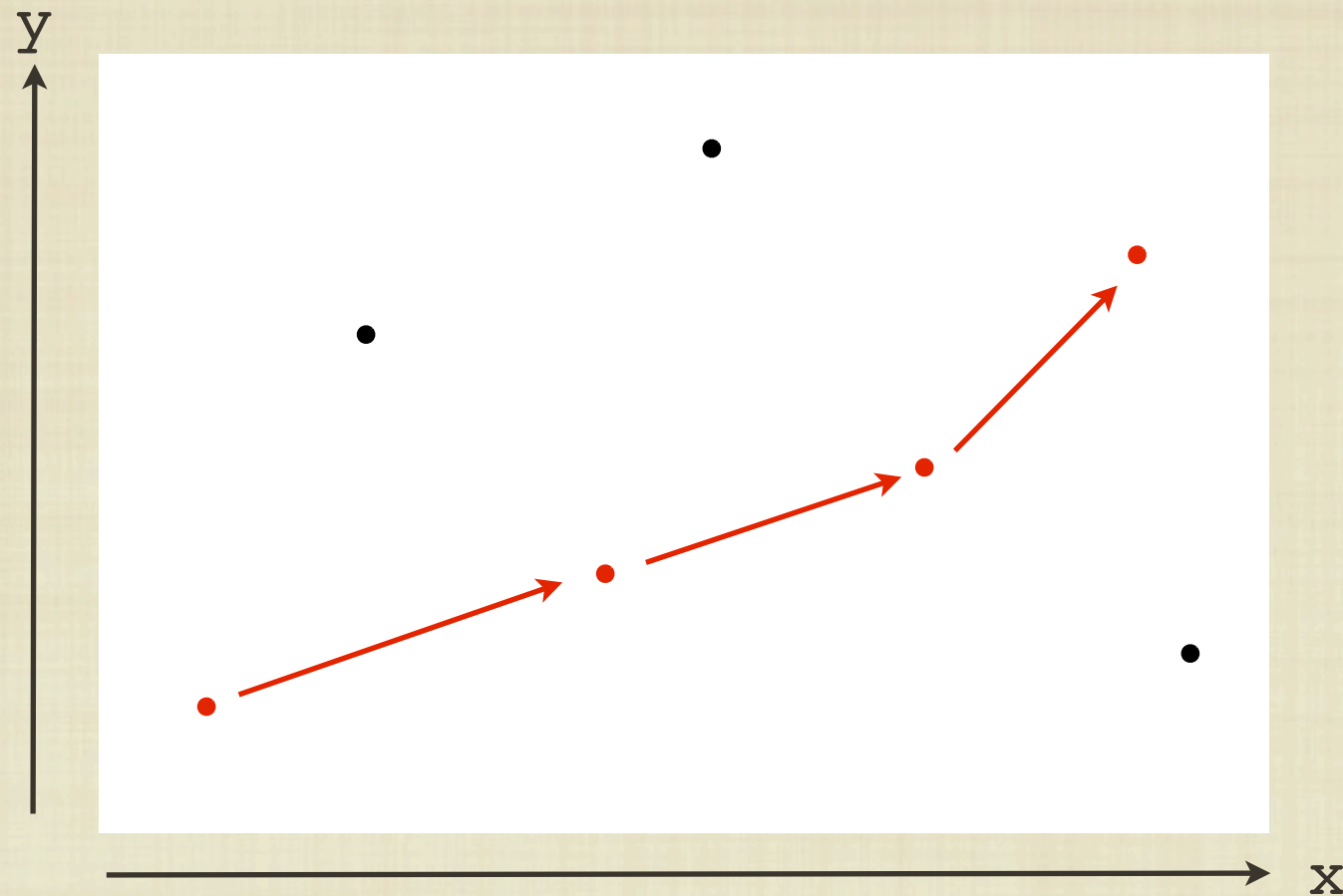
- When you want to plant a tree with height h at time t , you should plant a tree with height $h-t$ instead. Then you can ignore the growth of the trees.
- And plot trees on a x - y plane: x -coordinate is the position and y -coordinate is the modified height $(h-t)$.

SOLUTION



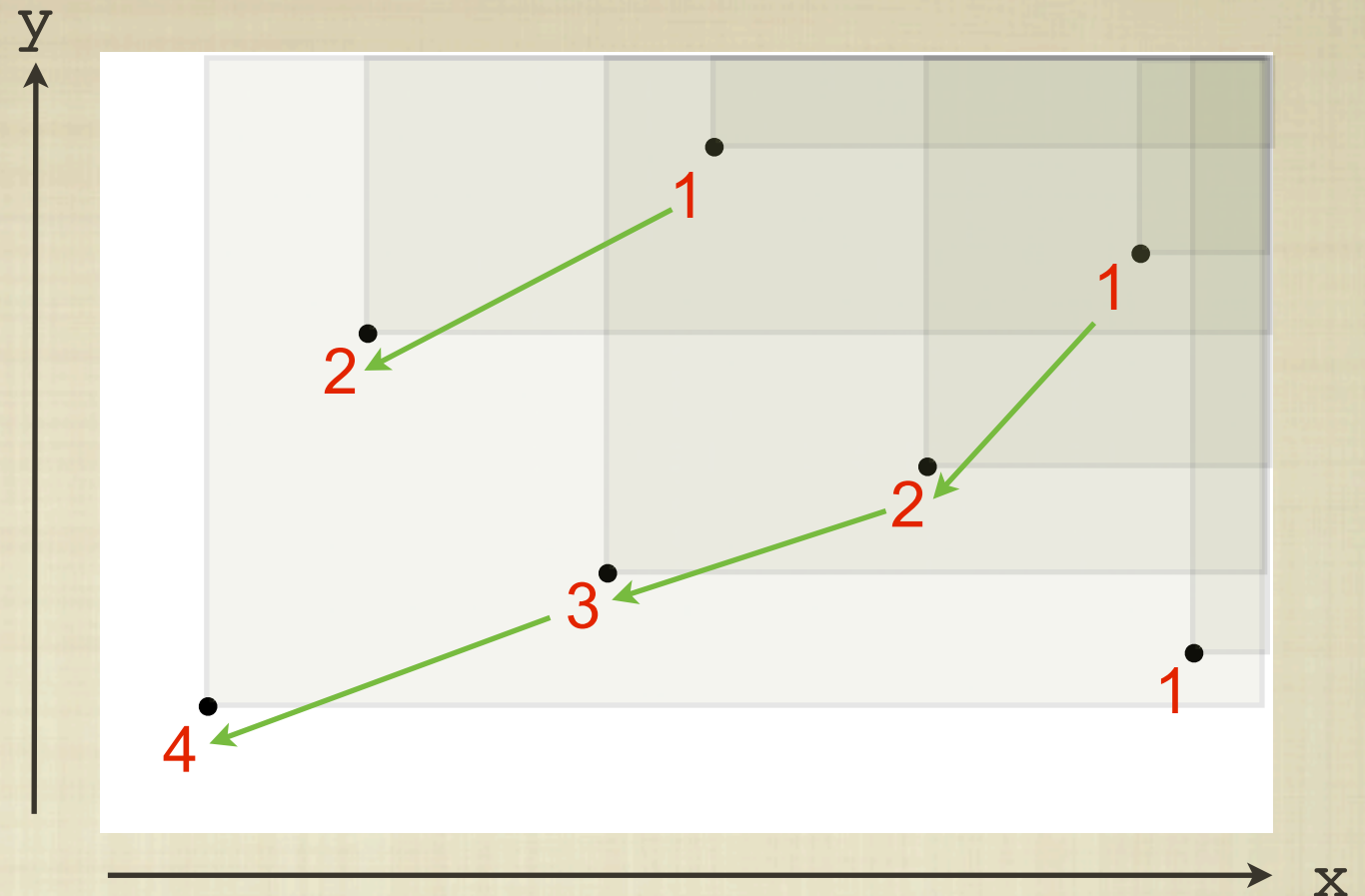
- like this picture
- then....

SOLUTION



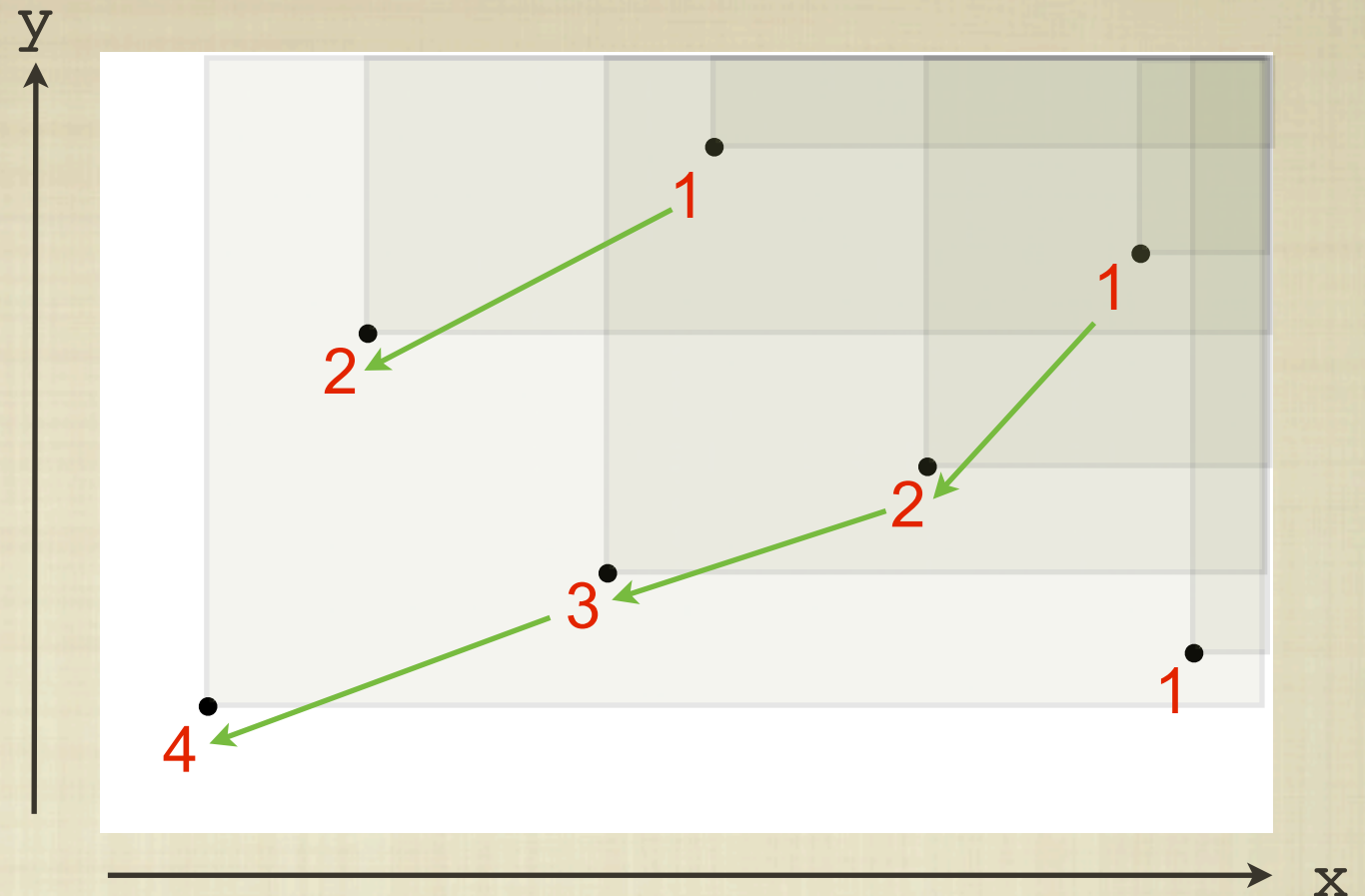
- LIS is the longest sequence of points P_1, P_2, \dots such that
- $x(P_1) < x(P_2) < \dots$
- $y(P_1) < y(P_2) < \dots$

SOLUTION



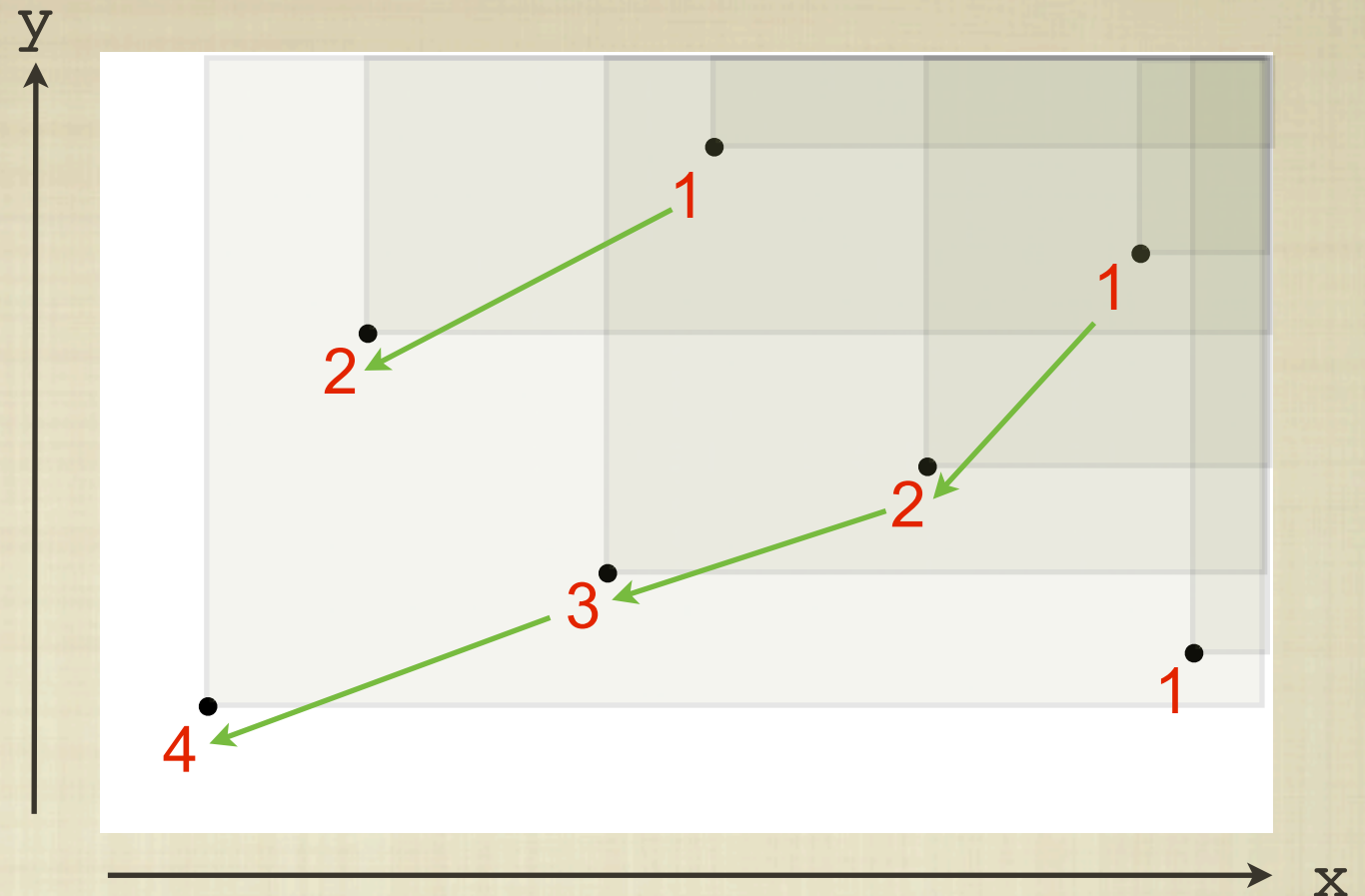
- At each point write the length of the longest increasing sequence that starts from the point.
- The value written on point $p =$
(the maximal value written in the rectangle whose lower-left corner is p) + 1.

SOLUTION



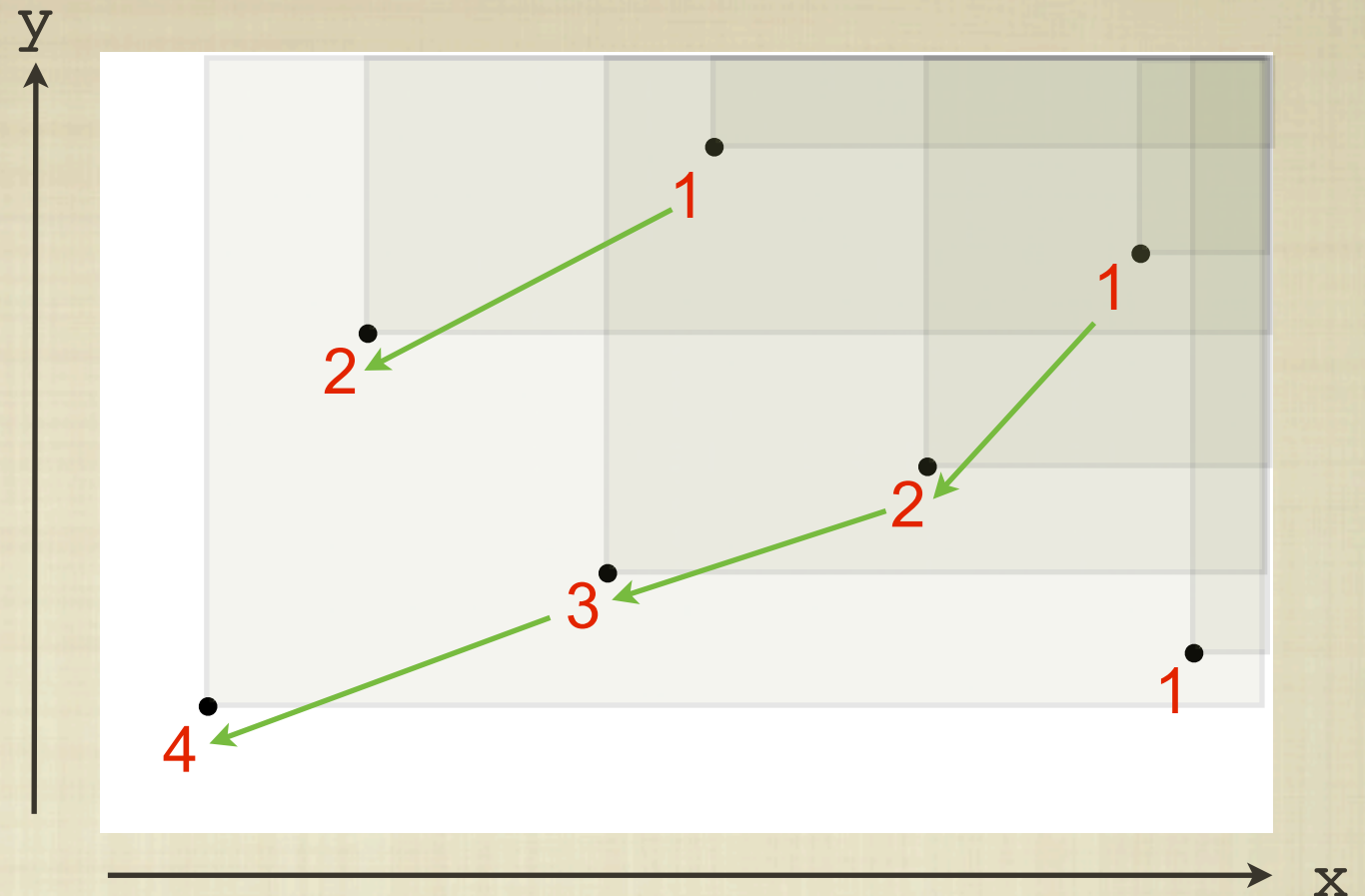
- How should we process queries?

SOLUTION



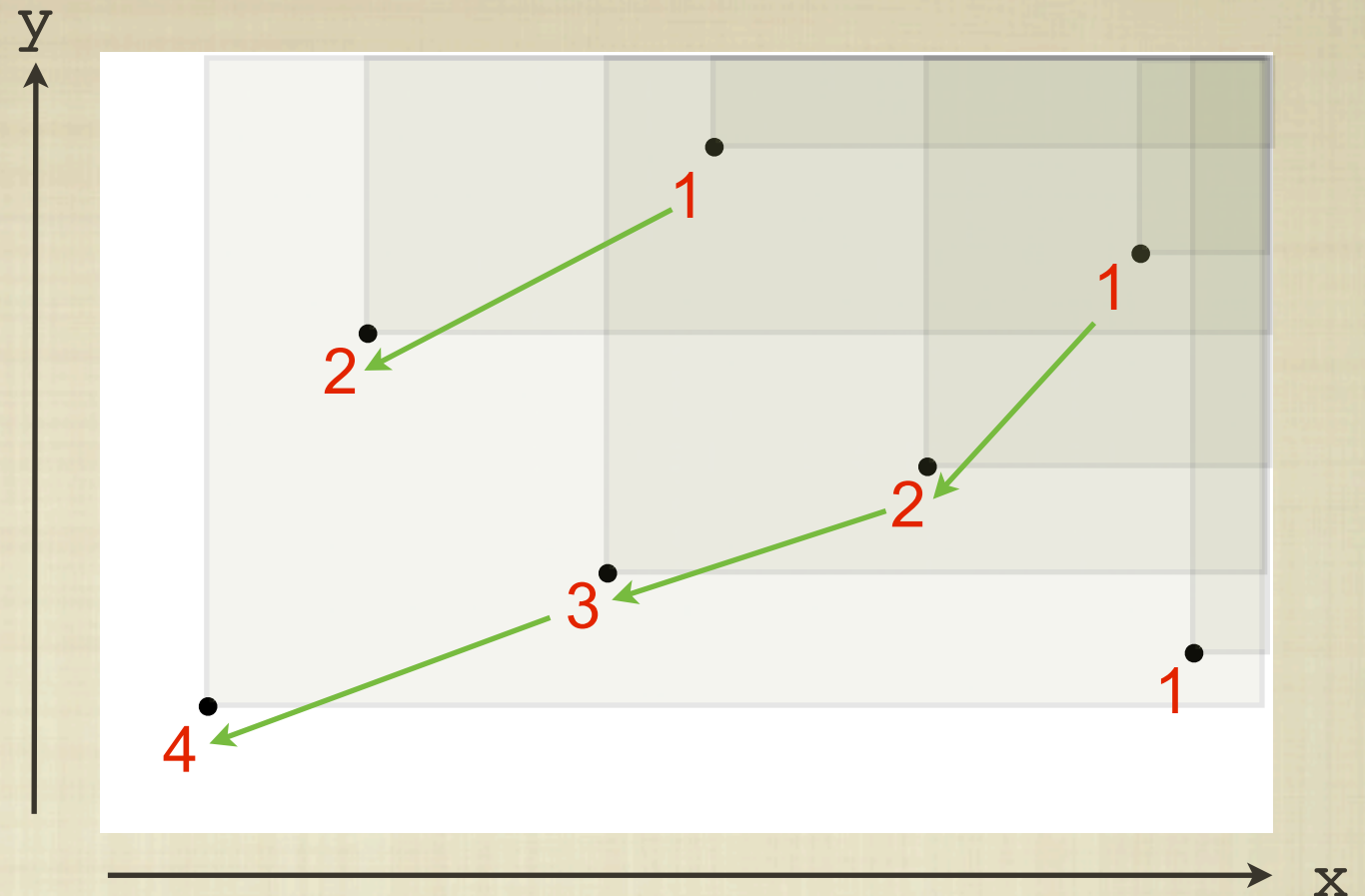
- Planting query \rightarrow plot a new point
- The new point will be one of the ten points that have the smallest y-coordinate.
- To process this query, erase all values written below the new point first and rewrite the values to those points from top to bottom.

SOLUTION



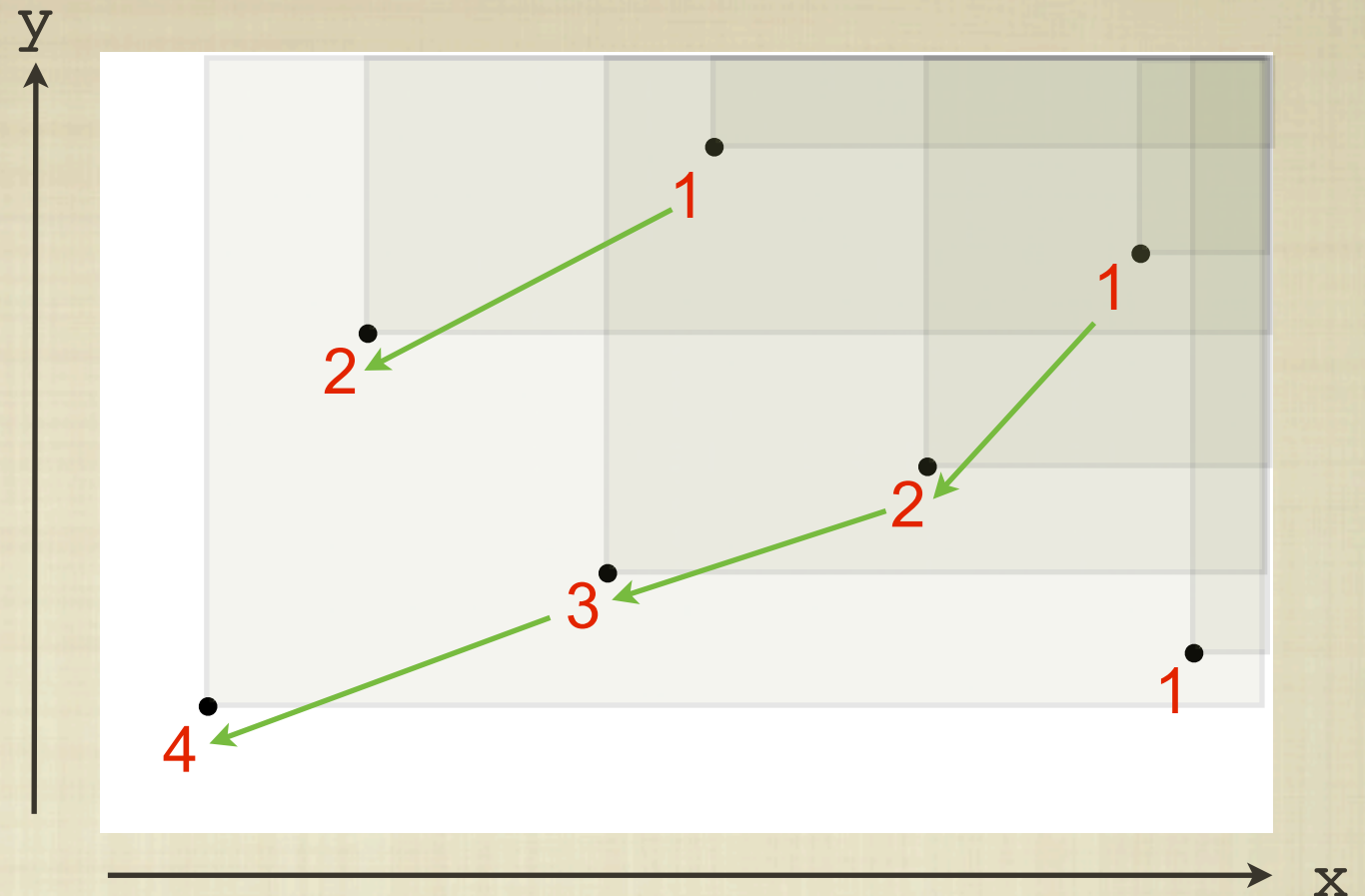
- Cutting query \rightarrow remove a point
- Similarly the removed point will be one of the ten points that have the smallest x-coordinate.
- So you can erase all values written on those points and rewrite correct values from right to left.

SOLUTION



- What data structures do we need?
- 2D segtree?
 - too slow :(

SOLUTION



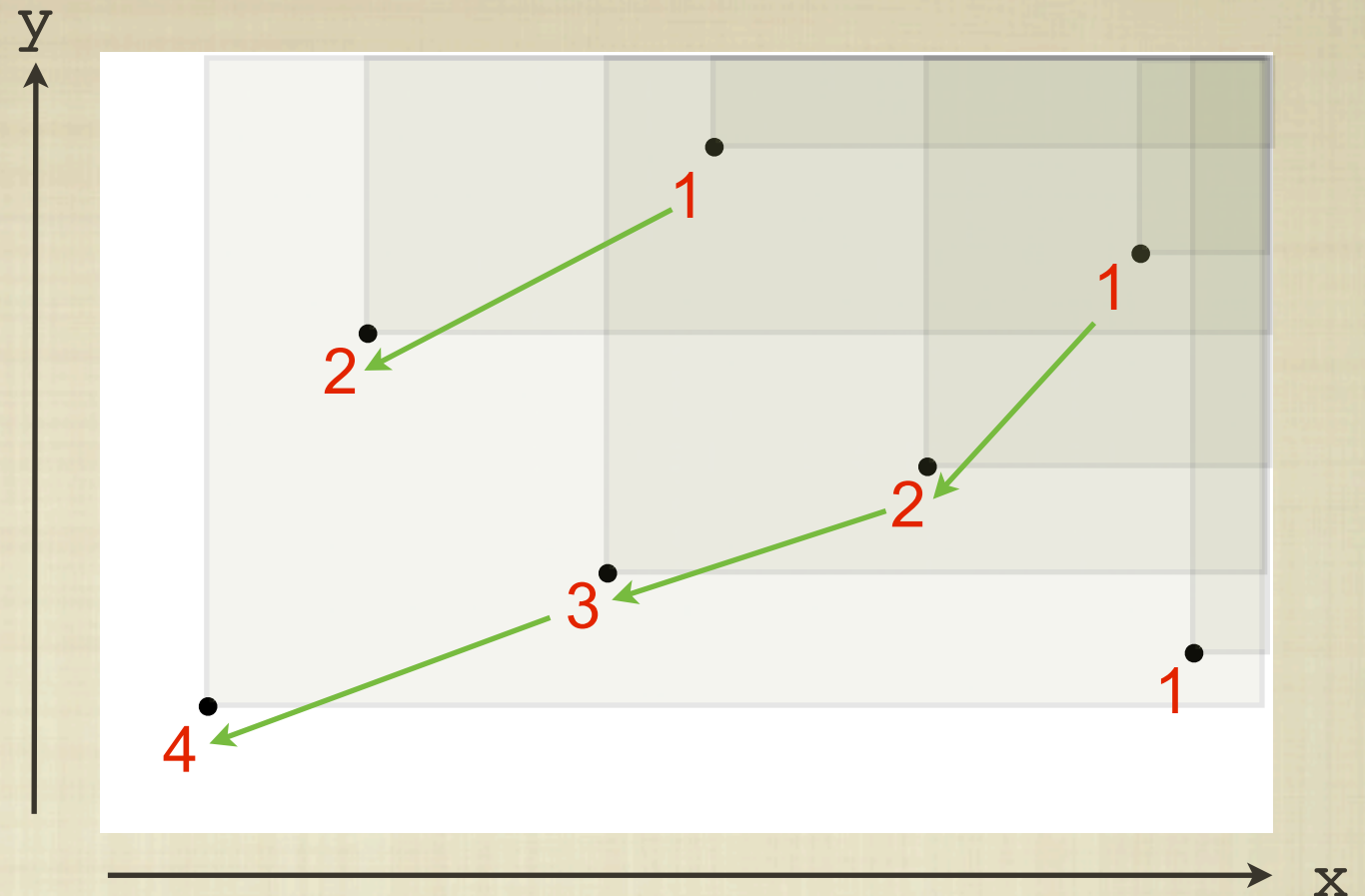
- Make two segment trees: let's call them `segx` and `segy`.

`segx` : x-directional segtree

`segy` : y-directional segtree

- For planting queries use `segx`.
- For cutting queries use `segy`.

SOLUTION



- The i -th leaf of seg_x contains the value written on the point whose x -coordinate is x and non-leaf nodes of the segment trees have the maximum of children of the node.
- Define seg_y similarly.

TIME COMPLEXITY

- $O(N * 10 * \log N)$

THANK YOU FOR WATCHING!