

THE UNIVERSITY OF MELBOURNE

DOCTORAL THESIS

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# The Coupling Time for the Ising Heat-Bath Dynamics & Efficient Optimization for Statistical Inference

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*Abstract*

Faculty of Science

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Doctor of Philosophy

**The Coupling Time for the Ising Heat-Bath Dynamics & Efficient  
Optimization for Statistical Inference**

by Timothy HYNDMAN

The title page must be followed by an abstract of 300–500 words in English. The Thesis Abstract is written here (and usually kept to just this page). The page is kept centered vertically so can expand into the blank space above the title too.

# Declaration of Authorship

This is to certify that:

1. the thesis comprises only my original work towards the PhD except where indicated in the Preface,
2. due acknowledgement has been made in the text to all other material used,
3. the thesis is fewer than 100 000 words in length, exclusive of tables, maps, bibliographies and appendices.

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# Preface

If applicable, a Preface page includes a statement of:

- Work carried out in collaboration indicating the nature and proportion of the contribution of others and in general terms the portions of the work which the candidate claims as original
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- Where a substantially unchanged multi-author paper is included in the thesis a statement prepared by the candidate explaining the contributions of all involved. A signed copy by all authors must be included with the submission form.

*“Thanks to my solid academic training, today I can write hundreds of words on virtually any topic without possessing a shred of information, which is how I got a good job in journalism.”*

Dave Barry

# *Acknowledgements*

The acknowledgements and the people to thank go here, don't forget to include your project advisor...

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*For/Dedicated to/To my...*

# Chapter 1

## Introduction to this Thesis

Initially, this thesis was intended to be made up entirely of the contents of Part [II](#), along with what we hoped would be several significant further contributions to the study. However, the practicalities of a deadline, along with the challenging nature of the research, meant that the decision was made to augment this thesis with an essentially separate section of study. This is what makes up Part [I](#).

The reader should view these two parts as standalone topics, to be read independently. However, they are not without any commonality. Both are within the realm of stochastic mathematics, Part [I](#) being a study of a random variable constructed from a stochastic process, and Part [II](#) being a study of probability distributions that maximize certain statistical objective functions.

## Part I

# The Coupling Time for the Ising Heat-Bath Dynamics

## Chapter 2

# Introduction

### 2.1 The Ising Model

The Ising model is named after Ernst Ising who studied it in his 1924 thesis [1] under the supervision of Wilhelm Lenz who invented the model [2]. It was originally motivated by the phenomenon of ferromagnetism but it has since been re-interpreted to apply to numerous other situations in both physics and other fields <sup>1</sup>.

The Ising model has enjoyed a prominent position in the statistical physics literature. This is largely due to the existence of a phase transition, a sharp transition in the large scale behaviour of the model as a parameter moves past a critical value. The transition was first shown to exist by Rudolph Peierls [4] in what was the first proof of the existence of a phase transition for any model in statistical mechanics. Additionally, the Ising model is both relatively simple, and also mathematically tractable in some non-trivial cases [5]. These qualities are rare among models with a phase transition and so the Ising model has become somewhat of a staple for both studying phase transitions and testing new statistical mechanics techniques.

The model is a probability distribution on spin configurations - assignments of +1 and -1 spins to each vertex in a finite graph  $G = (V, E)$ . The set of all possible configurations is

$$\Omega = \{-1, +1\}^V \tag{2.1}$$

and for a particular configuration,  $\sigma \in \Omega$ , we refer to the spin of a particular vertex  $i \in V$  with  $\sigma[i]$ . Each configuration has an associated energy, given by

$$H_{G,\beta,h}(\sigma) = -\beta \sum_{ij \in E} \sigma[i]\sigma[j] - h \sum_{i \in V} \sigma[i] \tag{2.2}$$

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<sup>1</sup>See [3, notes of Section 1.4.2] for a list of references concerning this.

where  $\beta \in [0, \infty)$  is the inverse temperature, and  $h \in \mathbb{R}$  is the magnetic field.

The Gibbs measure is the distribution on  $\Omega$  that characterises the Ising model and it is defined by

$$\pi_{G,\beta,h}(\sigma) \propto \exp(-H_{G,\beta,h}(\sigma)). \quad (2.3)$$

In everything that follows, we will be concerned only with the zero-field ( $h = 0$ ) Ising model. This gives us the slightly simpler form for the Gibbs measure,

$$\pi_{G,\beta}(\sigma) \propto \exp\left(\beta \sum_{ij \in E} \sigma[i]\sigma[j]\right), \quad \sigma \in \{-1, 1\}^V. \quad (2.4)$$

### 2.1.1 The Phase Transition

An in depth study of the Ising phase transition and its associated critical temperature will not be needed for this work. However, we will still wish to refer to it occasionally and so here we give a workable description of the phase transition on lattices.

Consider the Gibbs measure with zero-field (2.4) in the limits  $\beta \downarrow 0$  and  $\beta \uparrow \infty$ . It is easy to see that in the former limit, the measure is uniform across all configurations and in the latter limit, the measure assigns all weight to the constant configurations  $\sigma^- = (-1, -1, \dots, -1)$  and  $\sigma^+ = (+1, +1, \dots, +1)$ . This leads to the following overly simplistic description of the phase transition. It is the change in distribution that occurs as we increase the temperature; from distributions concentrated on states whose spins mostly agree, to distributions producing states which have roughly equal numbers of plus and minus spins.

To be slightly more concrete we define quantities called the magnetization and magnetization density. The *magnetization* on a volume  $\Lambda \subseteq V$  is defined as

$$M_\Lambda(\sigma) = \sum_{i \in \Lambda} \sigma[i]. \quad (2.5)$$

Normalizing this gives the *magnetization density*,  $M_\Lambda(\sigma)/|\Lambda|$ . On the  $d$ -dimensional torus with side length  $L$ ,  $G(L) = (\mathbb{Z}/L\mathbb{Z})^d$ , the quantity

$$L(\beta) = \lim_{L \rightarrow \infty} \mathbb{E}_\beta \left| \frac{M_{G(L)}(\sigma)}{|G(L)|} \right| \quad (2.6)$$

depends on the inverse temperature  $\beta$ . When  $d = 1$ ,  $L(\beta) = 0$  for any  $\beta$  and there is no phase transition. However, when  $d > 1$ , there exists some critical  $\beta_c(d)$  such that  $L(\beta) = 0$  for  $\beta < \beta_c(d)$  and  $L(\beta) > 0$  for  $\beta > \beta_c(d)$  [3]. This  $\beta_c(d)$  is the critical inverse temperature at which we observe a phase transition.

## 2.2 Coupling from the Past

One of the primary concerns regarding the Ising model is how to efficiently sample from the Gibbs measure. Calculating the normalizing constant for (2.4), known as the partition function, is a #P-complete problem [6]. As such a direct approach to sampling is computationally intractable and so other methods must be employed instead. One such method is Markov Chain Monte Carlo (MCMC). This involves constructing a Markov chain whose states are elements of  $\Omega$  and whose stationary distribution is given by (2.4). One can then obtain a sample by running this Markov chain for long enough that the output has distribution sufficiently close to (2.4). One difficulty in using MCMC is that, initially, one does not know how long to run the chain for. In principal, bounds on this time can be achieved, but in practise, proving these bounds can be very challenging.

An alternative to MCMC was introduced by Propp and Wilson called Coupling from the Past (CFTP) [7]. Unlike MCMC, CFTP not only has an automatically determined running time, but it has the additional advantage of outputting exact samples from the stationary distribution. This does not come without a cost - CFTP has a random running time. Therefore, a key question towards evaluating the effectiveness of CFTP is understanding the distribution of its running time, that is, the *coupling time*.

In Chapters 3 and 4, we will investigate the coupling time for the Ising heat-bath Glauber dynamics, both on the cycle, at any temperature, in Chapter 3, and on any vertex transitive graph, at sufficiently high temperatures, in Chapter 4. Our main result in each chapter will be proving that, when appropriately scaled, the coupling time essentially converges to a Gumbel distribution as the size of the graph increases.

### 2.2.1 Ising heat-bath Glauber dynamics

The continuous-time heat-bath Glauber dynamics for the Ising model is a Markov chain whose states are elements of  $\Omega$  and whose stationary distribution is given by (2.4). For a given graph  $G = (V, E)$ , and a given inverse temperature,  $\beta$ , we can describe the dynamics as follows.

Initialize every vertex in  $V$  with a spin (for example, we could start in the all-plus configuration). To each vertex in  $V$  we give an i.i.d. rate-one Poisson clock. Define the probability

$$p_i(\sigma) = \frac{e^{\beta S_i(\sigma)}}{e^{\beta S_i(\sigma)} + e^{-\beta S_i(\sigma)}} \quad (2.7)$$

where

$$S_i(\sigma) = \sum_{j \sim i} \sigma[j] \quad (2.8)$$

is the sum of the spins of the neighbours of  $i$ , and  $j \sim i$  denotes that  $j$  is connected to  $i$  with some edge  $ij \in E$ . Let  $\sigma_t$  denote the spin configuration at time  $t$ . When the clock of vertex  $i$  rings at some time  $t$ , we update  $\sigma_t[i]$  to  $+1$  with probability  $p_i(\sigma_t)$ , and to  $-1$  otherwise.

### 2.2.2 The Coupling Time

We now describe the two coupled chains from which we define the coupling time of the Ising heat-bath Glauber dynamics. In order to do this, it will prove convenient to use a random mapping representation for the jump process. This will also help us outline an implementation of CFTP for our coupling.

Define  $f : \Omega \times V \times [0, 1] \mapsto \Omega$  via  $f(\sigma, i, u) = \sigma'$  where  $\sigma'[j] = \sigma[j]$  for  $j \neq i$  and

$$\sigma'[i] = \begin{cases} 1, & u \leq p_i(\sigma), \\ -1, & u > p_i(\sigma). \end{cases} \quad (2.9)$$

Let  $\mathcal{V}$  and  $U$  be independent, with  $\mathcal{V}$  uniform on  $V$  and  $U$  uniform on  $[0, 1]$ . Then, updating our chain at rate  $n = |V|$ , and performing updates from  $\sigma$  to  $\sigma'$  according to  $\sigma' = f(\sigma, \mathcal{V}, U)$ , we recover the dynamics described in Section 2.2.1.

We also note that  $f$  is monotonic, in the following sense. We define a partial ordering on  $\Omega$  by writing that  $\sigma \preceq \omega$  if  $\sigma, \omega \in \Omega$  are such that  $\sigma[i] \leq \omega[i]$  for all  $i \in V$  (and similarly for  $\sigma \succeq \omega$ ). Then for any fixed  $i \in V$  and  $u \in [0, 1]$ , if  $\sigma \preceq \omega$  then  $f(\sigma, i, u) \preceq f(\omega, i, u)$ .

Let  $(\mathcal{V}_k, U_k)_{k \geq 1}$  be an i.i.d. sequence of copies of  $(\mathcal{V}, U)$ . Define top and bottom chains,  $(\mathcal{T}_t)_{t \geq 0}$  and  $(\mathcal{B}_t)_{t \geq 0}$ , with initial states

$$\mathcal{T}_0 = (1, 1, \dots, 1) \quad (2.10)$$

$$\mathcal{B}_0 = (-1, -1, \dots, -1) \quad (2.11)$$

that update together at rate  $n$ . On the  $k$ th update at time  $t_k$ , update  $\mathcal{T}_{t_k}$  to  $f(\mathcal{T}_{t_k}, \mathcal{V}_k, U_k)$  and update  $\mathcal{B}_{t_k}$  to  $f(\mathcal{B}_{t_k}, \mathcal{V}_k, U_k)$ .

We call the coupled process,  $(\mathcal{B}_t, \mathcal{T}_t)_{t \geq 0}$ , *the Ising heat-bath coupling*. From the monotonicity of  $f$ ,  $\mathcal{T}_t \succeq \mathcal{B}_t$ , for all  $t \geq 0$ .

A more descriptive explanation of the coupling is that the top and bottom chains share the same rate-one Poisson clocks at each vertex, and upon updating that vertex, we share



the same uniform random variable  $U$  between the two chains to determine whether to update to a plus or minus according to (2.9).

The *coupling time* of the Ising heat-bath process is the random variable

$$T = \inf \{t : \mathcal{T}_t = \mathcal{B}_t\}. \quad (2.12)$$

This is the main object of interest for our analysis. Note that the coupling time is not just a property of the Ising heat-bath process, but also of the coupling we have chosen. In Section 3.1 we will make a change to the coupling we use to make the analysis easier. Some care will need to be taken to verify that the coupling time is not affected by this change.

### 2.2.3 Equivalence of Discrete and Continuous Coupling Time

So far we have stated that the running time of CFTP has the same distribution as the coupling time. In fact, we have glossed over one important detail. Namely, CFTP is exclusively run in discrete time, and our coupling time is defined by the continuous time dynamics. Therefore, for our motivation to be reasonable, we would like to show some sort of equivalence between the distributions of the discrete and continuous coupling times. We do this via Claim 2.2 which first requires the following Lemma.

**Lemma 2.1.** *Let  $T(k)$  be the sum of  $k$  i.i.d. rate  $\lambda$  exponentials. For all  $0 < \epsilon < 1$ ,*

$$\mathbb{P} \left( \left| \frac{T(k)\lambda}{k} - 1 \right| \geq \epsilon \right) \leq 2 \exp(-k\epsilon^2/4) \quad (2.13)$$

*Proof.*

$$\mathbb{P} \left( \left| \frac{T(k)\lambda}{k} - 1 \right| \geq \epsilon \right) = \mathbb{P} \left( \frac{T(k)\lambda}{k} \leq 1 - \epsilon \right) + \mathbb{P} \left( \frac{T(k)\lambda}{k} \geq 1 + \epsilon \right) \quad (2.14)$$

Since  $T(k)$  is the sum of  $k$  i.i.d. rate  $\lambda$  exponentials, its moment generating function is

$$M_k(t) = \left( \frac{\lambda}{\lambda - t} \right)^k, \quad t < \lambda. \quad (2.15)$$

Using a Chernoff bound, for all  $0 < t < \lambda$ ,

$$\mathbb{P} \left( \frac{T(k)\lambda}{k} \geq 1 + \epsilon \right) = \mathbb{P} \left( T(k) \geq \frac{k}{\lambda}(1 + \epsilon) \right) \quad (2.16)$$

$$\leq \left( \frac{\lambda}{\lambda - t} \right)^k \exp \left( -\frac{tk}{\lambda}(1 + \epsilon) \right) \quad (2.17)$$

$$= \exp(k(\ln(\lambda/(\lambda - t)) - t(1 + \epsilon)/\lambda)). \quad (2.18)$$

Taking  $t = \lambda - \lambda/(1 + \epsilon)$ ,

$$\mathbb{P}\left(\frac{T(k)\lambda}{k} \geq 1 + \epsilon\right) \leq \exp(k(\log(1 + \epsilon) - \epsilon)). \quad (2.19)$$

Similarly, for all  $t < 0$ ,

$$\mathbb{P}\left(\frac{T(k)\lambda}{k} \leq 1 - \epsilon\right) = \mathbb{P}\left(T(k) \leq \frac{k}{\lambda}(1 - \epsilon)\right) \quad (2.20)$$

$$\leq \left(\frac{\lambda}{\lambda - t}\right)^k \exp\left(-\frac{tk}{\lambda}(1 - \epsilon)\right) \quad (2.21)$$

$$= \exp(k(\ln(\lambda/(\lambda - t)) - t(1 - \epsilon)/\lambda)). \quad (2.22)$$

Taking  $t = \lambda - \lambda/(1 - \epsilon)$ ,

$$\mathbb{P}\left(\frac{T(k)\lambda}{k} \leq 1 - \epsilon\right) \leq \exp(k(\log(1 - \epsilon) + \epsilon)). \quad (2.23)$$

Overall,

$$\mathbb{P}\left(\left|\frac{T(k)\lambda}{k} - 1\right| \geq \epsilon\right) \leq \exp(k(\log(1 - \epsilon) + \epsilon)) + \exp(k(\log(1 + \epsilon) - \epsilon)) \quad (2.24)$$

$$\leq 2 \exp(k(\log(1 + \epsilon) - \epsilon)) \quad (2.25)$$

$$\leq 2 \exp(-k\epsilon^2/4) \quad (2.26)$$

for  $0 < \epsilon < 1$ . □

**Claim 2.2.** Let  $(N_n)_{n \in \mathbb{N}}$  be a sequence of positive random integers, and  $(m_n)_{n \in \mathbb{N}}$  be a non-decreasing sequence of integers such that  $N_n \geq m_n$  for all  $n$  and  $\lim_{n \rightarrow \infty} m_n = \infty$ . Define  $T(n)$  be the random time it takes for a rate  $\lambda$  Poisson clock to go off  $n$  times (this is the time it takes for  $n$  updates to occur in a continuous-time Markov Chain).

Let  $a_n$  and  $b_n$  be positive deterministic sequences such that  $b_n/a_n \rightarrow \infty$  and for any  $c > 0$

$$\frac{m_n a_n^2}{b_n^2} - c \log \frac{b_n}{a_n} \rightarrow \infty. \quad (2.27)$$

Define

$$Y_n = \frac{T(N_n) - b_n}{a_n} \quad (2.28)$$

and

$$Z_n = \frac{N_n - \lambda b_n}{\lambda a_n}. \quad (2.29)$$

Let  $X$  be a random variable which doesn't place any mass at infinity. Then  $Y_n \xrightarrow{d} X$  if and only if  $Z_n \xrightarrow{d} X$ .

*Proof.* In order to prove either direction, it is sufficient to show that for any  $\epsilon > 0$ ,

$$\lim_{n \rightarrow \infty} \mathbb{P}(|Y_n - Z_n| > \epsilon) = 0. \quad (2.30)$$

First note that

$$|Y_n - Z_n| = \left| \frac{T(N_n) - b_n}{a_n} - \frac{N_n - \lambda b_n}{\lambda a_n} \right| \quad (2.31)$$

$$= \left| \frac{T(N_n)}{a_n} - \frac{N_n}{\lambda a_n} \right| \quad (2.32)$$

$$= \left| \frac{T(N_n)\lambda}{N_n} - 1 \right| \frac{N_n}{\lambda a_n}. \quad (2.33)$$

So for any  $\epsilon > 0$

$$\mathbb{P}(|Y_n - Z_n| > \epsilon) \leq \mathbb{P}\left(\left|\frac{T(N_n)\lambda}{N_n} - 1\right| > \epsilon \frac{a_n}{4b_n}\right) + \mathbb{P}\left(\frac{N_n}{\lambda a_n} > \frac{4b_n}{a_n}\right). \quad (2.34)$$

We will show that both of the terms on the right hand side vanish as  $n \rightarrow \infty$ . We start with the first of these.

Since  $N_n \geq m_n$ ,

$$\mathbb{P}\left(\left|\frac{T(N_n)\lambda}{N_n} - 1\right| > \epsilon \frac{a_n}{4b_n}\right) \leq \mathbb{P}\left(\sup_{k \geq m_n} \left|\frac{T(k)\lambda}{k} - 1\right| > \epsilon \frac{a_n}{4b_n}\right) \quad (2.35)$$

$$= \mathbb{P}\left(\bigcup_{k \geq m_n} \left\{\left|\frac{T(k)\lambda}{k} - 1\right| > \epsilon \frac{a_n}{4b_n}\right\}\right) \quad (2.36)$$

$$\leq \sum_{k=m_n}^{\infty} \mathbb{P}\left(\left|\frac{T(k)\lambda}{k} - 1\right| > \epsilon \frac{a_n}{4b_n}\right). \quad (2.37)$$

To apply Lemma 2.1, we need that  $\epsilon a_n/(4b_n) < 1$ . However, since  $a_n/b_n \rightarrow 0$ , we can ensure this holds by taking  $n$  large enough. Continuing,

$$\sum_{k=m_n}^{\infty} \mathbb{P}\left(\left|\frac{T(k)\lambda}{k} - 1\right| > \epsilon \frac{a_n}{4b_n}\right) \leq 2 \sum_{k=m_n}^{\infty} \exp\left(-k\epsilon^2 \frac{a_n^2}{64b_n^2}\right) \quad (2.38)$$

$$= 2 \frac{\exp(-\epsilon^2 a_n^2 (m_n - 1)/(64b_n^2))}{\exp(\epsilon^2 a_n^2/(64b_n^2)) - 1}. \quad (2.39)$$

Since  $x \leq \exp(x) - 1$  for  $0 \leq x \leq 1$ , for sufficiently large  $n$ ,

$$2 \frac{\exp(-\epsilon^2 a_n^2 (m_n - 1)/(64b_n^2))}{\exp(\epsilon^2 a_n^2/(64b_n^2)) - 1} \leq \frac{128b_n^2}{a_n^2 \epsilon^2} \exp(-\epsilon^2 a_n^2 (m_n - 1)/(64b_n^2)) \quad (2.40)$$

$$\leq 256 \frac{b_n^2}{a_n^2 \epsilon^2} \exp(-m_n \epsilon^2 a_n^2/(64b_n^2)) \quad (2.41)$$

By (2.27), this goes to zero as  $n \rightarrow \infty$ .

To bound the second term in (2.34), we will treat the two directions of the proof separately. Firstly, assume that  $Z_n \xrightarrow{d} X$ . Then note that

$$\mathbb{P}\left(\frac{N_n}{\lambda a_n} > \frac{4b_n}{a_n}\right) = \mathbb{P}\left(Z_n > 3\frac{b_n}{a_n}\right) \quad (2.42)$$

and since  $b_n/a_n \rightarrow \infty$ , and  $Z_n$  converges to a distribution which places no mass at  $\infty$ ,

$$\lim_{n \rightarrow \infty} \mathbb{P}\left(\frac{N_n}{\lambda a_n} > \frac{4b_n}{a_n}\right) = 0. \quad (2.43)$$

Now assume that  $Y_n \xrightarrow{d} X$ . Note that if  $T(N_n)/a_n < \epsilon/2$  and  $|T(N_n)\lambda/N_n - 1| < 1/2$ , then  $N_n/(\lambda a_n) < \epsilon$ . So

$$\mathbb{P}\left(\frac{N_n}{\lambda a_n} > \frac{4b_n}{a_n}\right) \leq \mathbb{P}\left(\frac{T(N_n)}{a_n} < \frac{2b_n}{a_n}\right) + \mathbb{P}\left(\left|\frac{T(N_n)\lambda}{N_n} - 1\right| > \frac{1}{2}\right) \quad (2.44)$$

$$= \mathbb{P}\left(Y_n < \frac{b_n}{a_n}\right) + \mathbb{P}\left(\left|\frac{T(N_n)\lambda}{N_n} - 1\right| > \frac{1}{2}\right). \quad (2.45)$$

As above, since  $a_n/b_n \rightarrow \infty$ , and  $Y_n$  converges to a distribution which places no mass at  $\infty$ , the first term vanishes in the limit. The second disappears since

$$\mathbb{P}\left(\left|\frac{T(N_n)\lambda}{N_n} - 1\right| > \frac{1}{2}\right) \leq \mathbb{P}\left(\left|\frac{T(N_n)\lambda}{N_n} - 1\right| > \epsilon \frac{a_n}{4b_n}\right) \quad (2.46)$$

for large enough  $n$ .

□

*Remark 2.3.* We apply Claim 2.2 to the coupling time of the Glauber heat-bath dynamics in the following way. Take  $N_n$  to be the discrete coupling time on a graph of size  $n$ . The continuous time coupling time is given by  $T(N_n)$ . Note that  $N_n \geq m_n = n$  since each vertex must be updated at least once for coupling to occur. Finally Theorems 3.1 and 4.1 establish the limiting distribution of the continuous-time coupling time using scaling and shifting sequences  $a_n$  and  $b_n$  whose ratio is

$$\frac{b_n}{a_n} = \log n \quad (2.47)$$

and thus (2.27) is satisfied. This means that, appropriately scaled, the discrete-time coupling time has the same limiting distribution as the continuous-time coupling time.

### 2.2.4 Summary of CFTP

We are now in a position to give a brief summary of the CFTP method, as it applies to the Ising heat-bath coupling. It should be noted that we include this summary of CFTP for completeness. None of the details regarding the implementation of CFTP are required outside of this section. It serves only as motivation for the study of the coupling time.

Let  $f : \Omega \times V \times [0, 1] \mapsto \Omega$  and  $(\mathcal{V}, U)$  be as defined in Section 2.2.2. Let  $(\mathcal{V}_k, U_k)$  be an i.i.d. sequence of copies of  $(\mathcal{V}, U)$  and define

$$f_{-k} = f(\cdot, \mathcal{V}_k, U_k). \quad (2.48)$$

We construct the composition

$$F_{-k} = f_0 \circ f_1 \circ \cdots \circ f_{k-1} \quad (2.49)$$

and define the *backwards coupling time* to be

$$T_{\text{BACK}} = \min\{k \in \mathbb{N} : F_{-k}(\mathcal{B}_0) = F_{-k}(\mathcal{T}_0)\}. \quad (2.50)$$

The state  $F_{-T_{\text{BACK}}}(\mathcal{B}_0) = F_{-T_{\text{BACK}}}(\mathcal{T}_0)$  is the output of the CFTP algorithm, and was shown by Propp and Wilson [7] to be an exact sample from the chain's stationary distribution. To see why this is so, observe that by the monotonicity of  $f$ , if  $F_{-k}(\mathcal{B}_0) = F_{-k}(\mathcal{T}_0)$ , then  $F_{-k}(\sigma) = F_{-k}(\mathcal{B}_0)$  for any  $\sigma \in \Omega$ . If we let  $\sigma_\pi$  be a random sample from the stationary distribution  $\pi$ , then  $F_{-k}(\mathcal{B}_0) = F_{-k}(\mathcal{T}_0) = F_{-k}(\sigma_\pi)$  must also have distribution  $\pi$ , which in our case is given by (2.4).

If we reverse the composition to construct

$$F_k = f_k \circ f_{k-1} \circ \cdots \circ f_1 \quad (2.51)$$

we can define the usual discrete time coupling time as

$$T_{\text{DIS}} = \min\{k \in \mathbb{N} : F_k(\mathcal{B}_0) = F_k(\mathcal{T}_0)\}. \quad (2.52)$$

The forwards coupling time,  $T_{\text{DIS}}$ , has the same distribution as the backwards coupling time,  $T_{\text{BACK}}$ , although in general,  $F_{T_{\text{DIS}}}(\mathcal{B}_0) = F_{T_{\text{DIS}}}(\mathcal{T}_0)$  does not have distribution (2.4).

In practise, one runs the CFTP algorithm by starting both the top and bottom chains from some point in the past to time zero. This is repeated for increasingly more distant

times in the past until both chains agree at time 0. The sequence of times at which one restarts this process need not be  $-1, -2, -3, \dots$ , rather, any monotonic natural sequence  $a_1, a_2, \dots$  can be used. See [8], [9], and [10] for further discussion.

## 2.3 Information percolation

A cornerstone to the proofs contained in Chapters 3 and 4 is the framework of information percolation, first introduced by Lubetzky and Sly in 2016 [11]. In this paper, Lubetzky and Sly managed to achieve much sharper results, in much more generality, regarding the mixing time for the Glauber dynamics for the Ising model than had been achieved before. In this section we provide a brief summary of their paper before laying out the basic framework, in the context of the Ising heat-bath dynamics, that will be required for Chapters 3 and 4.

### 2.3.1 Information percolation and cutoff for the stochastic Ising model

In order to define cutoff, the central phenomenon of study in Lubetzky and Sly's 2016 paper titled, 'Information percolation and cutoff for the stochastic Ising model', we first have to define the total-variation mixing time. Given a parameter  $\epsilon$ , a Markov Chain  $Y_t$  has mixing time

$$t_{\text{MIX}}(\epsilon) = \inf \left\{ t : \max_{x_0 \in \Omega} \|\mathbb{P}(X_t \in \cdot | X_0 = x_0) - \pi\|_{\text{TV}} \leq \epsilon \right\} \quad (2.53)$$

where the total variation distance  $\|\nu_1 - \nu_2\|_{\text{TV}}$  is defined as

$$\max_{A \in \Omega} |\nu_1(A) - \nu_2(A)| = \frac{1}{2} \sum_{\sigma \in \Omega} |\nu_1(\sigma) - \nu_2(\sigma)|. \quad (2.54)$$

A family of Markov chains  $(Y_t)$  indexed by  $n$  is said to exhibit cutoff if

$$t_{\text{MIX}}(\epsilon) = (1 + o(1))t_{\text{MIX}}(\epsilon'), \quad (2.55)$$

for any fixed  $0 < \epsilon, \epsilon' < 1$ . A *cutoff window* is a sequence  $w_n$  where

$$t_{\text{MIX}}(\epsilon) = t_{\text{MIX}}(1 - \epsilon) + \mathcal{O}(w_n) \quad (2.56)$$

for any  $0 < \epsilon < 1$ .

Historically, proving cutoff has proven to be highly challenging. In a survey on the topic, Diaconis [12] wrote 'proof of a cutoff is a difficult, delicate affair, requiring detailed

knowledge of the chain, such as all eigenvalues and eigenvectors'. It is therefore worth noting the significant gap between the strength of the results regarding cutoff achieved using information percolation, and those that existed previously.

Previous to [11], the best result known for general graphs was that cutoff occurs with a  $\mathcal{O}(1)$  window in the simple case when  $\beta = 0$  [13]. However, no results were known for  $\beta > 0$ , despite a conjecture by Levin et al. in 2009 [8, Section 23.2] that cutoff occurs on any sequence of transitive graphs when the mixing time is of order  $\log n$  (as one would expect when  $\beta < c_0$  for some  $c_0 > 0$  that depends on the sequence of graphs). On lattices, the first results to appear were due to Lubetzky and Sly in 2013 who established cutoff up to the critical temperature for dimensions  $d \leq 2$  with only a  $\mathcal{O}(\log \log n)$  window [14].

Using information percolation, Lubetzky and Sly proved the existence of cutoff for the continuous time Glauber dynamics for the Ising model with an  $\mathcal{O}(1)$  window on  $\mathbb{Z}^d$  for all temperatures up to the critical temperature. In a companion paper [15], they extended this result to include any graph for sufficiently high temperature. Since these papers, information percolation has also been used to establish cutoff for the Swendsen-Wang dynamics on the lattice [16], suggesting that the technique is effective on a broader class of problems than simply Glauber dynamics for Ising.

### 2.3.2 The framework

At its core, information percolation is a way of tracking how the dependencies of the final spins of the Glauber heat-bath dynamics percolate through the graph over time. These dependencies are traced backwards through time from some designated time  $t^*$  on the space-time slab  $V \times [0, t^*]$  (see Figure 2.1 for example) to create the update history. These histories are made in such a way so that, if for any  $j \in V$  no path exists connecting  $(i, t^*)$  to  $(j, 0)$ , then the spin of  $i$  does not depend on the initial state (and thus at time  $t^*$  vertex  $i$  takes  $+1$  and  $-1$  spins with equal probability by symmetry). The main constructs used to create this history are the update sequence, and the update support function which we will now define.

#### 2.3.2.1 The update sequence

Recalling our random mapping representation from Section 2.2.2, we can encode an update of our coupled process with the tuple  $(\mathcal{V}, U, t)$ , where  $t$  is the time of the update,  $\mathcal{V}$  is the vertex that is updated, and  $U$  is the value of the uniform random variable that tells us whether  $\mathcal{V}$  is a plus or minus according to (2.9). The *update sequence* along an



**Figure 2.1** – A section of the space-time slab  $V \times [0, t^*]$  along with a typical appearance of the update histories for two vertices on the cycle. Time runs vertically from bottom to top, and the vertices are represented by circles, laid out horizontally. If there is a path in the update history of  $v$  between points  $(u, t)$  and  $(v, t^*)$ , then the spin of  $v$  at time  $t^*$  depends on the spin of  $u$  at time  $t$ . In this example, since there is no path from vertex  $(i, t^*)$  to time 0, the final spin at  $i$  does not depend on the initial configuration whereas the final spin at  $j$  does.

interval  $(t_0, t_1]$  is the set of these tuples with  $t_0 < t \leq t_1$ . Given the state of our Markov Chain at time  $t_0$ ,  $Y_{t_0}$ , the update sequence along  $(t_0, t_1]$  contains all the information we need to construct  $Y_{t_1}$ . In particular, given the update sequence along the interval  $(0, t_1]$ ,  $Y_{t_1}$  is a deterministic function of  $Y_0$ .

### 2.3.2.2 The update support function

Given the update sequence along the interval  $(t_1, t_2]$ , the *update support function*,  $\mathcal{F}(A, t_1, t_2)$ , is the minimal set of vertices whose spins at time  $t_1$  determine the spins of the vertices in  $A$  at time  $t_2$ . That is,  $i \in \mathcal{F}(A, t_1, t_2)$  if and only if there exist states  $Y_{t_1}, Y'_{t_1} \in \{-1, +1\}^V$  that differ only at  $i$  and such that when we construct  $Y_{t_2}$  and  $Y'_{t_2}$  using the update sequence,  $Y_{t_2} \neq Y'_{t_2}$ .

In particular, if  $\mathcal{F}(i, 0, t) = \emptyset$  then the spin at vertex  $i$  at time  $t$  does not depend on the initial state and so for any two coupled chains  $Y$  and  $Y'$ ,  $Y_t[i] = Y'_t[i]$ . As a consequence of the monotonicity of our coupling, we can make the stronger statement



that  $\mathcal{T}_t[i] = \mathcal{B}_t[i]$  if and only if  $\mathcal{F}(i, 0, t) = \emptyset$  which of course means that

$$\mathbb{P}[\mathcal{T}_t[i] \neq \mathcal{B}_t[i]] = \mathbb{P}[\mathcal{F}(i, 0, t) \neq \emptyset]. \quad (2.57)$$

For ease of notation, we will often use the shorthand

$$\mathcal{H}_i(t) := \mathcal{F}(i, t, t^*) \quad (2.58)$$

where  $t^*$  is some target time that should be clear from context. We call this the *update history of vertex  $i$  at time  $t$* . Tracing  $\mathcal{H}_i(t)$  backwards in time from  $t^*$  produces a subgraph of  $\Omega \times [0, t^*]$  which we write as  $\mathcal{H}_i$  and which we simply call the *update history of vertex  $i$* . To be slightly more precise, to produce  $\mathcal{H}_i$  we connect  $(j, t)$  to  $(j, t')$  if  $j \in \mathcal{H}_i(t)$  and there are no updates of  $j$  along  $(t', t]$  and we connect  $(j, t)$  to  $(j', t)$  if there was an update at  $(j, t)$ ,  $j \in \mathcal{H}_i(t)$ ,  $j' \notin \mathcal{H}_i(t)$ , and  $j' \in \mathcal{H}_i(t + \epsilon)$  for any sufficiently small  $\epsilon > 0$ .

Similarly, we also use

$$\mathcal{H}_A(t) := \mathcal{F}(A, t, t^*) \quad (2.59)$$

for the update history of a vertex set  $A$  at time  $t$  and  $\mathcal{H}_A$  for the update history of vertex set  $A$ . Note that

$$\mathcal{H}_A(t) = \bigcup_{i \in A} \mathcal{H}_i(t) \quad (2.60)$$

and

$$\mathcal{H}_A = \bigcup_{i \in A} \mathcal{H}_i. \quad (2.61)$$

To give some intuition to the definitions above, consider how we might construct the update history of a vertex  $i$  from some target time  $t^*$ . We have at our disposal the update sequence along  $(0, t^*]$  which we place in order of decreasing time. If vertex  $i$  does not appear in the update sequence then we create a temporal edge between  $(i, t^*)$  and  $(i, 0)$  and our update history is complete - vertex  $i$  was never updated and so it simply takes its initial value. Otherwise, we create temporal edge between  $(i, t^*)$  and  $(i, t_i)$  where  $t_i$  is the last time vertex  $i$  was updated. At this point we note from (2.9) that the spin that vertex  $i$  takes due to this update depends on the spins of its neighbours. So we add spatial edges from  $(i, t_i)$  to  $(j, t_i)$  for each  $j \sim i$ . Finally, we can iterate this process for each neighbour until every history has reached time 0.

In Figure 2.2 we have followed this procedure to show how we might create the update history from a single vertex on the cycle. This construction certainly contains every vertex that might influence the final spin at  $i$ , that is, it contains  $\mathcal{H}_i$  as a subgraph.



**Figure 2.2** – A naive construction of the update history of  $i$ . Each update  $(\mathcal{V}, U, t)$  in the update sequence is represented by a  $*$  at  $(\mathcal{V}, t)$ .

However, it is possible for updates to occur that do not depend on neighbouring spins. These updates cause temporal edges leading up to them to terminate without branching out to the neighbouring vertices. These type of updates are called *oblivious updates*.

### 2.3.2.3 Oblivious updates

Generally speaking, an update to a vertex is oblivious if we do not need to know the configuration of its neighbours to determine the spin of that vertex. More precisely, an update,  $(\mathcal{V}, U, t)$ , is oblivious if

$$f(\sigma, \mathcal{V}, U)[\mathcal{V}] = f(\sigma', \mathcal{V}, U)[\mathcal{V}] \quad (2.62)$$

for all  $\sigma, \sigma' \in \Omega$ , where  $f$  is as defined in (2.9).

Consider how these updates occur under our random mapping representation. Let  $\Delta_i$  denote the degree of a vertex  $i$ . Recalling (2.7),

$$\frac{e^{-\beta\Delta_i}}{e^{\beta\Delta_i} + e^{-\beta\Delta_i}} \leq p_i(\sigma) \leq \frac{e^{\beta\Delta_i}}{e^{\beta\Delta_i} + e^{-\beta\Delta_i}}, \quad (2.63)$$

with equality holding for the lower and upper limits when the neighbours have spins all minus and all plus respectively. So for a particular update  $(\mathcal{V}, U, t)$ , if  $U \leq \frac{e^{-\beta\Delta_{\mathcal{V}}}}{e^{\beta\Delta_{\mathcal{V}}} + e^{-\beta\Delta_{\mathcal{V}}}}$  then  $\mathcal{V}$  is updated to a plus regardless of the configuration of its neighbours. Hence  $(\mathcal{V}, U, t)$  is an oblivious update. Similarly, if  $U > \frac{e^{\beta\Delta_{\mathcal{V}}}}{e^{\beta\Delta_{\mathcal{V}}} + e^{-\beta\Delta_{\mathcal{V}}}}$  then  $\mathcal{V}$  is updated to a

minus regardless of the configuration of its neighbours and hence  $(\mathcal{V}, U, t)$  is an oblivious update. It is easy to see that these are the only types of oblivious updates.

The rate of these updates at vertex  $i$  is

$$\theta_i = 1 - \left( \frac{e^{\beta\Delta_i}}{e^{\beta\Delta_i} + e^{-\beta\Delta_i}} - \frac{e^{-\beta\Delta_i}}{e^{\beta\Delta_i} + e^{-\beta\Delta_i}} \right) \quad (2.64)$$

$$= 1 - \tanh(\beta\Delta_i). \quad (2.65)$$

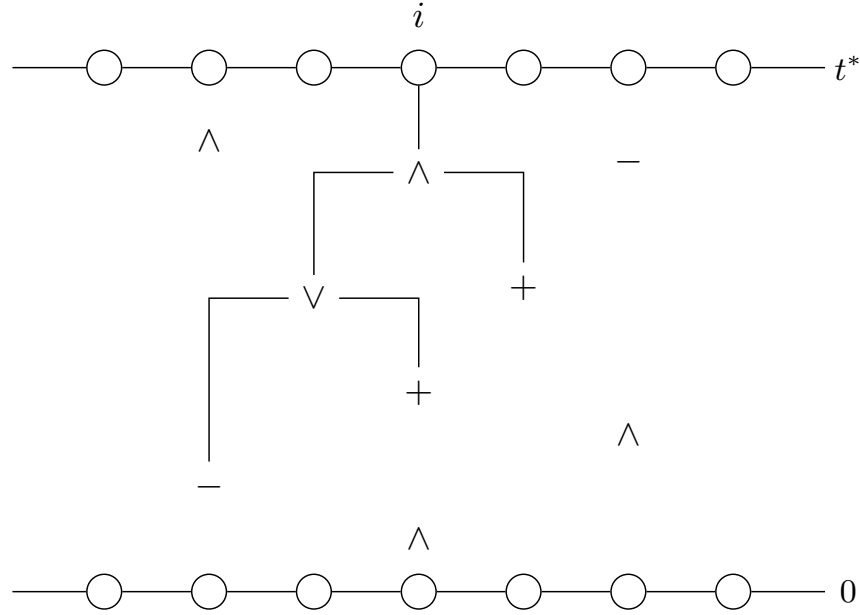
If  $G$  is a  $\Delta$ -regular graph (as will be the case in the following chapters) then we can drop the subscript and write  $\theta = 1 - \tanh(\beta\Delta)$  for the rate of oblivious updates at each vertex.

As noted earlier, oblivious updates cause temporal edges leading up to them in the update history to terminate. If  $j \in \mathcal{H}_i(t)$ , then an oblivious update  $(j, u, t)$  removes  $j$  from  $\mathcal{H}_i(t)$  without adding any of its neighbours. In Figure 2.3 we construct an update history from a single vertex  $i$  as in Figure 2.2, but this time terminate branches at oblivious updates. To help us represent the updates in our update sequence more precisely, note that on the cycle, the function defined in (2.9) can be rewritten as

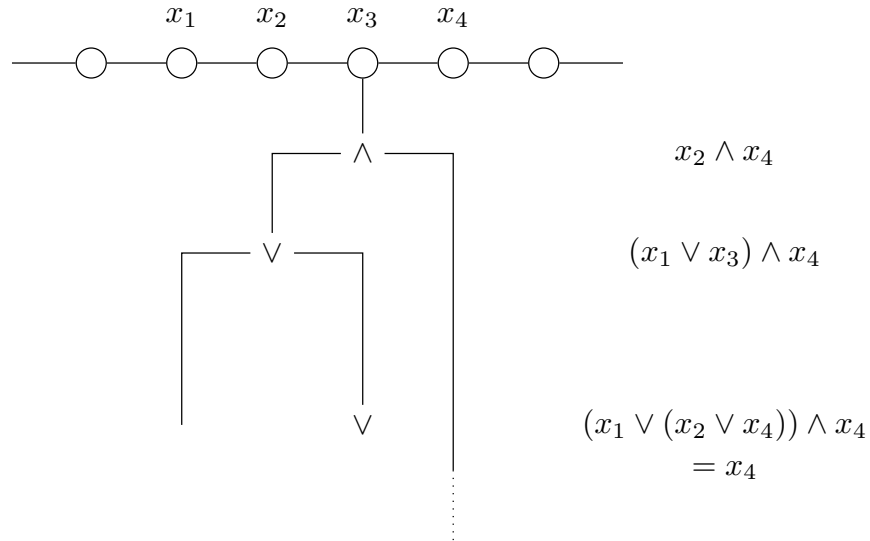
$$\sigma'[i] = \begin{cases} 1 & U \leq \theta/2, \\ \sigma[i-1] \vee \sigma[i+1] & \theta/2 < U \leq 1/2, \\ \sigma[i-1] \wedge \sigma[i+1] & 1/2 < U \leq 1 - \theta/2, \\ -1 & U > \theta/2. \end{cases} \quad (2.66)$$

We can therefore represent each update  $(\mathcal{V}, U, t)$  in the update sequence by placing at  $(\mathcal{V}, t)$  one of the symbols  $+$ ,  $\vee$ ,  $\wedge$ , or  $-$  chosen according to  $U$ . We then trace back from time  $t^*$ , branching to either side when we encounter a  $\vee$  or  $\wedge$ , and terminating whenever we encounter a  $+$  or  $-$ .

It is worth remarking that oblivious updates are not necessarily the only updates that can shrink the size of the update history of  $i$  (see for example Figure 2.4). However, for our analysis they will be the only such updates we will be concerned with. Indeed, in Chapter 3 we will use a different coupling so that these are the only updates that shrink the size of the update history, and in Chapter 4 we will use an alternative construction that bounds the true update history, in which updates are either oblivious or branch out to all  $\Delta$  neighbours.



**Figure 2.3** – The update sequence for a section of the cycle and the corresponding update history from vertex  $i$ . For this particular update sequence,  $i$  takes a final spin of  $+1$  regardless of the initial configuration.



**Figure 2.4** – A non-oblivious update that shrinks the size of the update history. On the right is written the final spin of  $x_3$  as a function of the configuration at that time. The update  $x_3 \mapsto x_2 \vee x_4$  causes the entire function to collapse to  $x_4$ , and so removes  $x_1$  and  $x_3$  from the update history.

## Chapter 3

# The Coupling Time on the Cycle

In this chapter we consider the Ising heat-bath Glauber dynamics (as described in Section 2.2.1) on the cycle  $G_n = (\mathbb{Z}/n\mathbb{Z})$ . The object of interest is the coupling time,  $T_n$ , which was defined in Section 2.2.2 but whose definition will be modified slightly in Section 3.1 to allow for simpler analysis. The main result is Theorem 3.1 which establishes that  $T_n$  converges in distribution to a Gumbel distribution at all temperatures. This confirms, for  $d = 1$ , a conjecture by Collevocchio et al. that the coupling time of the Ising heat-bath process on the lattice  $G_L = (\mathbb{Z}/L\mathbb{Z})^d$  converges to a Gumbel distribution as  $L \rightarrow \infty$  for all  $\beta < \beta_C$  [17, Conjecture 7.1] (We treat higher dimensions, and more generally any vertex transitive graphs, in Chapter 4). Of course, in one dimension, all temperatures are part of the high temperature regime [CITE], and likewise our result holds for any inverse-temperature  $\beta$ .

There is some intuition behind why we might expect that the coupling time converges to a Gumbel distribution. It is based on the belief that when the temperature is in the high-temperature regime, the dynamics behave qualitatively as if  $\beta = 0$ . In the  $\beta = 0$  case, the spins update independently of their neighbours, and thus the top and bottom chains can be coupled so that they agree on each vertex that has been updated. The coupling time is then precisely the time it takes for each vertex to be updated. This corresponds to the coupon collector's problem, which is known to have a Gumbel limit [18].

As mentioned in Section 2.2.4, the coupling time is of practical interest since its distribution is the same as that of the running time of the coupling from the past (CFTP) algorithm. Our result shows that when running the Glauber heat-bath dynamics for the Ising model on a large enough cycle, the running time of CFTP can be approximated by a Gumbel distribution. We note that even though one is typically more interested in the Ising model on lattices of dimension at least two (so that there exists a phase transition),

the one dimensional case proves to be a useful test case for the proof techniques. Furthermore, the applicability of Theorem 3.1 to the full high temperature regime justifies a separate treatment to the higher dimensional case in Chapter 4.

**Theorem 3.1.** *Let  $T_n$  be the coupling time for the continuous-time Ising heat-bath Glauber dynamics for the zero-field ferromagnetic Ising model on the cycle  $(\mathbb{Z}/n\mathbb{Z})$ . Then for any inverse-temperature  $\beta$ ,*

$$\lim_{n \rightarrow \infty} \mathbb{P} \left[ T_n < \frac{z + \ln n}{\theta} \right] = e^{-C_\theta e^{-z}} \quad (3.1)$$

where  $\theta = 1 - \tanh(2\beta)$  and  $C_\theta$  is a positive constant bounded by

$$\frac{1}{2\sqrt{\frac{4}{\theta} - 1} - 1} \leq \lambda \leq 1. \quad (3.2)$$

The proof of Theorem 3.1 will be given in Section 3.3 after the essential preliminaries are presented. In Section 3.1 we describe some modified dynamics and show that the coupling time we construct from these has the same distribution as the coupling time defined in Section 2.2.2. Then in Section 3.2 we outline the overall approach to the proof and define some essential quantities. Finally, Section 3.4 contains additional lemmas that are used in Section 3.3.

### 3.1 A new coupling on the cycle

On the cycle, we will use a different coupling of  $\mathcal{T}_t$  and  $\mathcal{B}_t$  via a new set of update rules that will replace those from (2.9). The new update rules simplify our update histories greatly by ensuring that each of the update histories never contain more than one vertex at any one time. However, we must be cautious. The coupling time is not just a property of the heat-bath dynamics, but also of the specific coupling we chose. Hence, we will have to verify that switching to our new rules does not change the distribution of  $T_n$ .

The new update rules are defined by using almost the same construction as in Section 2.2.2. The one difference is that we replace (2.9) as follows. When vertex  $i$  updates, instead of comparing  $U$  to the probability  $p_i(\sigma)$  to determine the new spin, we instead

$\mathbb{P}[(\mathcal{T}_t[i]', \mathcal{B}_t[i]') = \cdot]$		(1,1)	(1,-1)	(-1,-1)
$\mathcal{T}_t = \cdot$	$\mathcal{B}_t = \cdot$			
$(\dots, 1, \sigma_i, 1, \dots)$	$(\dots, 1, \sigma_i, 1, \dots)$	$1 - \theta$	0	$\frac{\theta}{2}$
$(\dots, 1, \sigma_i, 1, \dots)$	$(\dots, 1, \sigma_i, -1, \dots)$	$\frac{1}{2}$	$\frac{1-\theta}{2}$	$\frac{\theta}{2}$
$(\dots, 1, \sigma_i, 1, \dots)$	$(\dots, -1, \sigma_i, 1, \dots)$	$\frac{1}{2}$	$\frac{1-\theta}{2}$	$\frac{\theta}{2}$
$(\dots, 1, \sigma_i, 1, \dots)$	$(\dots, -1, \sigma_i, -1, \dots)$	$\frac{\theta}{2}$	$1 - \theta$	$\frac{\theta}{2}$
$(\dots, 1, \sigma_i, -1, \dots)$	$(\dots, 1, \sigma_i, -1, \dots)$	$\frac{1}{2}$	0	$\frac{1}{2}$
$(\dots, 1, \sigma_i, -1, \dots)$	$(\dots, -1, \sigma_i, -1, \dots)$	$\frac{\theta}{2}$	$\frac{1-\theta}{2}$	$\frac{1}{2}$
$(\dots, -1, \sigma_i, 1, \dots)$	$(\dots, -1, \sigma_i, 1, \dots)$	$\frac{1}{2}$	0	$\frac{1}{2}$
$(\dots, -1, \sigma_i, 1, \dots)$	$(\dots, -1, \sigma_i, -1, \dots)$	$\frac{\theta}{2}$	$\frac{1-\theta}{2}$	$\frac{1}{2}$
$(\dots, -1, \sigma_i, -1, \dots)$	$(\dots, -1, \sigma_i, -1, \dots)$	$\frac{\theta}{2}$	0	$1 - \theta$

**Table 3.1** – Probabilities of updating from  $(\mathcal{T}_t, \mathcal{B}_t)$  to  $(\mathcal{T}_t', \mathcal{B}_t')$  given vertex  $i$  updates at time  $t$ .

chose a new spin  $\sigma_i'$  via

$$\sigma_i' = \begin{cases} +1 & U < \theta/2, \\ \sigma_{i-1} & \theta/2 \leq U < 1/2, \\ \sigma_{i+1} & 1/2 \leq U < 1 - \theta/2, \\ -1 & U \geq 1 - \theta/2. \end{cases} \quad (3.3)$$

where  $U \in [0, 1]$  is an independent uniform random variable as before. It is easy to see that these update rules give rise to the same transition rates as those in (2.9). To show that the coupling time is unchanged, it is sufficient to verify that the joint jump probabilities of  $(\mathcal{T}_t[i], \mathcal{B}_t[i])$  are unchanged for each possible configuration of spins of vertices  $i - 1$  and  $i + 1$ . There are only nine possible configurations for the two neighbours of  $i$  in the top and bottom chain since  $\mathcal{B}_t[i] \leq \mathcal{T}_t[i], \forall t$ . Likewise, there are only three possible configurations for the updated spins  $(\mathcal{T}_t[i]', \mathcal{B}_t[i]')$ . Hence, given vertex  $i$  updates at time  $t$ , we can easily calculate all the required jump probabilities as shown in Table 3.1. These are unchanged whether using (2.9) or (3.3) and so the new rules do not change the coupled dynamics.





since our histories die at rate  $\theta$ .

### 3.2 Problem Setup

In order to prove Theorem 3.1, we will actually prove a stronger statement using Theorem 3.2. The general idea is that at some fixed time  $t^*$  we will count the number of vertices at which the bottom and top chains differ. This number is a random variable, which we will call  $W$ , and we can bound the total variation distance of its distribution with that of an appropriate compound poisson distribution. As a special case, we can use this bound as a bound on the probability that  $W$  is zero. Of course, if  $W$  is zero then the top and bottom chains must have coupled and so we can use this to establish Theorem 3.1.

Bounding the total variation distance between  $W$  and the compound Poisson will be done using compound Poisson approximation as described in [19]. This paper reviews a number of different methods by which approximations may be made. The specific method that we will employ is based on Stein's method for the compound Poisson distribution, introduced in [20].

We now make precise the ideas stated above. Fix  $z$  and a time of interest,  $t_* = (z + \ln n)/\theta$ . For each vertex  $i \in V$ , define indicators

$$X_i = \begin{cases} 1 & \mathcal{B}_{t_*}[i] \neq \mathcal{T}_{t_*}[i], \\ 0 & \mathcal{B}_{t_*}[i] = \mathcal{T}_{t_*}[i] \end{cases} \quad (3.5)$$

and set  $W = \sum_{i \in V} X_i$ . Note that from (3.4) we get

$$\mathbb{P}[X_i = 1] = e^{-\theta t_*} = \frac{e^{-z}}{n}. \quad (3.6)$$

For each  $i \in V$ , decompose  $W$  into  $W = X_i + U_i + Z_i + W_i$  where

$$U_i = \sum_{j \in B_i} X_j, \quad Z_i = \sum_{j \in C_i} X_j, \quad W_i = \sum_{j \in D_i} X_j. \quad (3.7)$$

and  $B_i, C_i$ , and  $D_i$  are the vertex sets

$$B_i = \{j \neq i : |j - i| \leq b_n\}, \quad (3.8)$$

$$C_i = \{j \notin B_i \cup \{i\} : |j - i| \leq c_n\}, \quad (3.9)$$

$$D_i = V \setminus (B_i \cup C_i \cup \{i\}). \quad (3.10)$$

We have some freedom in choosing  $b_n$  and  $c_n$ ; they are chosen to control the asymptotics of various quantities to be defined later. For this chapter, we will choose  $b_n = \ln(n)$  and  $c_n = \ln(n)^2$ .

We now define the quantities

$$\lambda = \sum_{i \in V} \mathbb{E} \left[ \frac{X_i}{X_i + U_i} I[X_i + U_i \geq 1] \right], \quad (3.11)$$

$$\mu_l = \frac{1}{l\lambda} \sum_{i \in V} \mathbb{E} [X_i I[X_i + U_i = l]], \quad l \geq 1, \quad (3.12)$$

which will be the parameters of the approximating compound Poisson distribution to  $W$ . We also define

$$\delta_1 = \sum_{i \in V} \sum_{k \geq 0} \mathbb{P}[X_i = 1, U_i = k] \mathbb{E} \left| \frac{\mathbb{P}[X_i = 1, U_i = k | W_i]}{\mathbb{P}[X_i = 1, U_i = k]} - 1 \right|, \quad (3.13)$$

$$\delta_4 = \sum_{i \in V} (\mathbb{E}[X_i Z_i] + \mathbb{E}[X_i] \mathbb{E}[X_i + U_i + Z_i]), \quad (3.14)$$

which we desire to be small for the compound Poisson approximation to be good.

The following theorem (reworked from [19]) bounds the distance between the distributions of  $W$  and the approximating compound Poisson.

**Theorem 3.2** ([19]). *Let  $W$ ,  $\lambda$ ,  $\boldsymbol{\mu}$ ,  $\delta_1$  and  $\delta_4$  be as defined above. Then*

$$d_{\text{TV}}(\mathcal{L}(W), \text{CP}(\lambda, \boldsymbol{\mu})) \leq (\delta_1 + \delta_4) e^\lambda. \quad (3.15)$$

Note that  $W$  is zero precisely when  $T < t^*$  and so the events  $\{W = 0\}$  and  $\{T \leq t^*\}$  are the same. Furthermore, if  $W'$  is compound Poisson with parameters  $\lambda$  and  $\boldsymbol{\mu}$ , where  $\boldsymbol{\mu}$  is supported only on the positive integers (as in (3.12)), then  $\mathbb{P}[W' = 0] = e^{-\lambda}$ . These observations lead to the following corollary of Theorem 3.2.

**Corollary 3.3.** *Let  $T_n$  be the coupling time of the continuous-time heat-bath Glauber dynamics for the zero-field Ising model at inverse-temperature  $\beta$  on the cycle  $(\mathbb{Z}/n\mathbb{Z})$  and let  $\delta_1$ ,  $\delta_4$  and  $\lambda$  be as defined above. Then*

$$\left| \mathbb{P} \left[ T_n \leq \frac{z + \ln(n)}{\theta} \right] - e^{-\lambda} \right| \leq (\delta_1 + \delta_4) e^\lambda, \quad (3.16)$$

where  $\theta = 1 - \tanh(2\beta)$ .

### 3.3 Proof of Theorem 3.1

In this section we use Corollary 3.3 to prove Theorem 3.1 by bounding  $\lambda$  and showing that  $\delta_1$  and  $\delta_4$  go to zero as  $n \rightarrow \infty$ . This is done in Lemmas 3.4, 3.5, and 3.6. The proofs of these require some additional lemmas concerning properties of the update histories which have been deferred to Section 3.4.

We begin by bounding  $\lambda$ .

**Lemma 3.4.** *Using the above setup*

$$\lim_{n \rightarrow \infty} \lambda = C_\theta e^{-z} \quad (3.17)$$

where

$$\frac{1}{2\sqrt{\frac{4}{\theta} - 1} - 1} \leq C_\theta \leq 1. \quad (3.18)$$

*Proof.* Beginning with the definition of  $\lambda$ , we have

$$\lambda = \sum_{i \in V} \mathbb{E} \left[ \frac{X_i}{X_i + U_i} I[X_i + U_i \geq 1] \right] \quad (3.19)$$

$$= \sum_{i=1}^n \mathbb{P}(X_i = 1) \mathbb{E} \left[ \frac{1}{1 + U_i} | X_i = 1 \right] \quad (3.20)$$

$$= \sum_{i=1}^n \frac{e^{-z}}{n} \mathbb{E} \left[ \frac{1}{1 + U_i} | X_i = 1 \right] \quad (3.21)$$

$$= e^{-z} \mathbb{E} \left[ \frac{1}{1 + U_i} | X_i = 1 \right] \quad (3.22)$$

where we have used that  $X_i$  is either zero or one, (3.6), and the transitivity of the graph in each step respectively. Clearly

$$\mathbb{E} \left[ \frac{1}{1 + U_i} | X_i = 1 \right] \leq 1. \quad (3.23)$$

By Jensen's inequality

$$\mathbb{E} \left[ \frac{1}{1 + U_i} | X_i = 1 \right] \geq \frac{1}{\mathbb{E}[1 + U_i | X_i = 1]} \quad (3.24)$$

$$= \frac{1}{1 + \mathbb{E}[U_i | X_i = 1]}. \quad (3.25)$$

Now

$$\mathbb{E}[U_i | X_i = 1] = \sum_{j \in B_i} \mathbb{P}[X_j = 1 | X_i = 1] \quad (3.26)$$

$$= \sum_{k=1}^{\lfloor b_n \rfloor} \sum_{|j-i|=k} \mathbb{P}[X_j = 1 | X_i = 1] \quad (3.27)$$

$$= 2 \sum_{k=1}^{\lfloor b_n \rfloor} \mathbb{P}[X_{i+k} = 1 | X_i = 1] \quad (3.28)$$

where we have used the symmetry of  $X_{i+k}$  and  $X_{i-k}$  in the last step. From Lemma 3.9,

$$\mathbb{E}[U_i | X_i = 1] \leq 2 \sum_{k=1}^{\lfloor b_n \rfloor} \left( \frac{e^{-z}}{n} + 2 \left( \frac{2 - \sqrt{\theta(4-\theta)}}{2-\theta} \right)^k \right) \quad (3.29)$$

$$< 2 \sum_{k=1}^{\lfloor b_n \rfloor} \frac{e^{-z}}{n} + 4 \sum_{k=1}^{\infty} \left( \frac{2 - \sqrt{\theta(4-\theta)}}{2-\theta} \right)^k \quad (3.30)$$

$$= \frac{2\lfloor b_n \rfloor}{n} e^{-z} + 2 \left( \sqrt{\frac{4}{\theta}} - 1 - 1 \right). \quad (3.31)$$

Finally, as  $n \rightarrow \infty$  the first term vanishes and

$$\liminf_{n \rightarrow \infty} \mathbb{E} \left[ \frac{1}{1 + U_i} | X_i = 1 \right] \geq \frac{1}{2\sqrt{\frac{4}{\theta}} - 1 - 1}. \quad (3.32)$$

To complete the proof, we need to show that the limit of

$$\mathbb{E} \left[ \frac{1}{1 + U_i} | X_i = 1 \right] \quad (3.33)$$

as  $n \rightarrow \infty$  exists. Since we already have bounds, it is sufficient to show that it is monotonic.

[NEED TO SHOW THAT THING ACTUALLY CONVERGES]

□

The next two lemmas prove that  $\delta_1$  and  $\delta_4$  go to zero as  $n \rightarrow \infty$ . Since from Lemma 3.4 we know that  $\lambda$  is bounded above by a constant, this is enough for (3.16) to show that

$$\lim_{n \rightarrow \infty} \mathbb{P} \left[ T_n < \frac{z + \ln n}{\theta} \right] = e^{-\lambda}. \quad (3.34)$$

**Lemma 3.5.** *Let  $\delta_1$  be as defined above in (3.13). Then*

$$\lim_{n \rightarrow \infty} \delta_1 = 0. \quad (3.35)$$

*Proof.* Starting with the definition of  $\delta_1$ , we have

$$\delta_1 = \sum_{i=1}^n \sum_{k=0}^{2\lfloor b_n \rfloor} \mathbb{P}[X_i = 1, U_i = k] \mathbb{E} \left| \frac{\mathbb{P}[X_i = 1, U_i = k | W_i]}{\mathbb{P}[X_i = 1, U_i = k]} - 1 \right|, \quad (3.36)$$

$$= n \sum_{k=0}^{2\lfloor b_n \rfloor} \mathbb{E} |\mathbb{P}[X_i = 1, U_i = k | W_i] - \mathbb{P}[X_i = 1, U_i = k]| \quad (3.37)$$

by the transitivity of the cycle. Let

$$C_i^c = \{j : |j - i| \leq (c_n + b_n)/2\} \quad (3.38)$$

be the set of vertices within distance  $(b_n + c_n)/2$  of  $i$  and define the events

$$A_1 = \{\exists j \in B_i \cup \{i\}, \exists t \in [0, t^*] : \mathcal{H}_j(t) \not\subseteq C_i^c\} \quad (3.39)$$

and

$$A_2 = \{\exists j \in D_i, \exists t \in [0, t^*] : \mathcal{H}_j(t) \cap C_i^c \neq \emptyset\} \quad (3.40)$$

as well as their intersection

$$A = A_1 \cap A_2. \quad (3.41)$$

From Lemma 3.7,

$$\mathbb{P}[X_i = 1, U_i = j | A^c, W_i] = \mathbb{P}[X_i = 1, U_i = j | A^c]. \quad (3.42)$$

Continuing on from (3.37), we split the probabilities into

$$\delta_1 = n \sum_{k=0}^{2\lfloor b_n \rfloor} \mathbb{E} \left| \mathbb{P}[X_i = 1, U_i = k | W_i, A] \mathbb{P}[A | W_i] - \mathbb{P}[X_i = 1, U_i = k | A] \mathbb{P}[A] + \right. \quad (3.43)$$

$$\left. \mathbb{P}(X_i = 1, U_i = k | A^c) (\mathbb{P}[A^c | W_i] - \mathbb{P}[A^c]) \right|$$

$$\leq n(2\lfloor b_n \rfloor + 1) \mathbb{E} \left[ \mathbb{P}[A | W_i] + \mathbb{P}[A] + \left| \mathbb{P}[A^c | W_i] - \mathbb{P}[A^c] \right| \right] \quad (3.44)$$

$$= n(2\lfloor b_n \rfloor + 1) \mathbb{E} [\mathbb{P}[A | W_i] + \mathbb{P}[A] + |1 - \mathbb{P}[A | W_i] - (1 - \mathbb{P}[A])|] \quad (3.45)$$

$$\leq n(2\lfloor b_n \rfloor + 1) \mathbb{E} [\mathbb{P}[A | W_i] + \mathbb{P}[A] + \mathbb{P}[A | W_i] + \mathbb{P}[A]] \quad (3.46)$$

$$= 2n(2\lfloor b_n \rfloor + 1) (\mathbb{E}[\mathbb{P}[A | W_i]] + \mathbb{P}[A]) \quad (3.47)$$

$$= 4n(2\lfloor b_n \rfloor + 1) \mathbb{P}[A]. \quad (3.48)$$

For either  $A_1$  or  $A_2$  to hold, there must exist a history that spreads at least distance  $(c_n - b_n)/2$  away from its starting vertex. By a union bound

$$\mathbb{P}[A] \leq \sum_{j=1}^n \mathbb{P}[\mathcal{H}_i \not\subseteq B(i, (c_n - b_n)/2) \times [0, t^*]] \quad (3.49)$$

$$= n \mathbb{P} \left[ \bigcup_{u \in [0, t^*]} \mathcal{H}_i(t^* - u) \not\subseteq B(i, (c_n - b_n)/2) \right] \quad (3.50)$$

Combining this with Lemma 3.8, and recalling our choices of  $b_n = \ln(n)$  and  $c_n = \ln(n)^2$  we get that

$$\delta_1 \leq 8n^2(2\lfloor b_n \rfloor + 1) \exp((z + \ln n)/\theta - \ln 2(c_n - b_n)/2) \quad (3.51)$$

$$\leq 16 \exp(z/\theta) n^{3+1/\theta + \ln 2/2 - \ln n \ln 2/2} \quad (3.52)$$

which goes to 0 as  $n \rightarrow \infty$ .  $\square$

**Lemma 3.6.** *Let  $\delta_4$  be as defined above in (3.14). Then*

$$\lim_{n \rightarrow \infty} \delta_4 = 0. \quad (3.53)$$

*Proof.* Starting with the definition of  $\delta_4$ , we have

$$\delta_4 = \sum_{i=1}^n (\mathbb{E}[X_i Z_i] + \mathbb{E}[X_i] \mathbb{E}[X_i + U_i + Z_i]), \quad (3.54)$$

$$= n \mathbb{P}[X_i = 1] \mathbb{E}[Z_i | X_i = 1] + e^{-z} \sum_{j \in \{i\} \cup B_i \cup C_i} \mathbb{E}[X_j], \quad (3.55)$$

$$= e^{-z} \mathbb{E}[Z_i | X_i = 1] + \frac{(2\lfloor c_n \rfloor + 1)e^{-2z}}{n}. \quad (3.56)$$

Now

$$\mathbb{E}[Z_i | X_i = 1] = \sum_{j \in C_i} \mathbb{P}[X_j = 1 | X_i = 1], \quad (3.57)$$

$$= 2 \sum_{k=\lfloor b_n \rfloor + 1}^{\lfloor c_n \rfloor} \mathbb{P}[X_{i+k} = 1 | X_i = 1]. \quad (3.58)$$

From Lemma 3.9,

$$\mathbb{E}[Z_i | X_i = 1] \leq 2 \sum_{k=\lfloor b_n \rfloor + 1}^{\lfloor c_n \rfloor} \left( \frac{e^{-z}}{n} + 2 \left( \frac{2 - \sqrt{\theta(4 - \theta)}}{2 - \theta} \right)^k \right), \quad (3.59)$$

$$\leq \frac{2(c_n - b_n + 1)e^{-z}}{n} + 4(c_n - b_n + 1) \left( \frac{2 - \sqrt{\theta(4 - \theta)}}{2 - \theta} \right)^{b_n + 1}. \quad (3.60)$$

Altogether,

$$\delta_4 \leq \frac{2(c_n - b_n + 1)e^{-z}}{n} + 4(c_n - b_n + 1) \left( \frac{2 - \sqrt{\theta(4 - \theta)}}{2 - \theta} \right)^{b_n + 1} + \frac{(2c_n + 1)e^{-2z}}{n} \quad (3.61)$$

which, recalling that  $b_n = \ln(n)$  and  $c_n = \ln(n)^2$ , goes to 0 as  $n \rightarrow \infty$ .  $\square$

### 3.4 Additional Lemmas

This section contains the proofs for a number of lemmas concerning properties of the update histories on the cycle.

**Lemma 3.7.** *Let  $i$  be a vertex in a graph and let*

$$C_i^c = \{j : |j - i| \leq (b_n - c_n)/2\} \quad (3.62)$$

*be the set of vertices within distance  $(b_n + c_n)/2$  of  $i$ . Define the events*

$$A_1 = \{\exists j \in B_i \cup \{i\}, \exists t \in [0, t^*] : \mathcal{H}_j(t) \not\subseteq C_i^c\} \quad (3.63)$$

*and*

$$A_2 = \{\exists j \in D_i, \exists t \in [0, t^*] : \mathcal{H}_j(t) \cap C_i^c \neq \emptyset\} \quad (3.64)$$

*as well as their union*

$$A = A_1 \cup A_2. \quad (3.65)$$

*Then*

$$\mathbb{P}[X_i = 1, U_i = j | A^c, W_i] = \mathbb{P}[X_i = 1, U_i = j | A^c]. \quad (3.66)$$

*Proof.* If  $A_1^c$  holds, then the events  $\{X_1 = 1\}$  and  $\{U_i = j\}$  depend only on the values of the update sequence inside  $C_i^c$ . If  $A_2^c$  holds then the events  $\{W_i = k\}$ ,  $k \geq 0$ , depend only on the values of the update sequence outside of  $C_i^c$ . Since the update sequences of each vertex are independent of each other vertex, if  $A_2^c$  holds, conditioning on  $W_i$  does

not affect the update sequences inside  $C_i^c$  and so

$$\mathbb{P}[X_i = 1, U_i = j | A^{\mathbb{G}}, W_i] = \mathbb{P}[X_i = 1, U_i = j | A^{\mathbb{G}}]. \quad (3.67)$$

[WOULD BE NICE TO SAY THIS MORE MATHEMATICALLY.]  $\square$

The following Lemma bounds how fast updates can percolate through the cycle. The proof is a slight modification of a similar one in [11].

**Lemma 3.8.** *Consider the update histories on the cycle. Let  $B(i, l)$  indicate the set of vertices at integer distance  $l$  or smaller from vertex  $i$ . The probability that the history of vertex  $i$  escapes  $B(i, l)$  in time  $s$  is bounded by*

$$\mathbb{P} \left[ \bigcup_{u \in [0, s]} \mathcal{H}_i(t^* - u) \not\subseteq B(i, l) \right] \leq 2 \exp(s - l \ln 2). \quad (3.68)$$

*Proof.* Let  $\mathbf{w}^- = (i, i-1, \dots, i-l)$  and  $\mathbf{w}^+ = (i, i+1, \dots, i+l)$  denote the sequences of adjacent vertices starting at vertex  $i$  and extending distance  $l$  to the left and right respectively. For  $\mathcal{H}_i$  to contain any vertex outside  $B(i, l)$  at a time  $u \in [t^* - s, t^*]$  then either each  $w_k^-$  was updated at some time  $t^* > t_k \geq t^* - s$  with  $t_{k-1} > t_k$  or each  $w_k^+$  was updated at some time  $t^* > t_k \geq t^* - s$  with  $t_{k-1} > t_k$ . Call the first event  $M_-$  and the second  $M_+$ . We have

$$\mathbb{P}[M_-] = \mathbb{P}[M_+] = \mathbb{P}[\text{Po}(s) \geq l] \quad (3.69)$$

where  $\text{Po}(s)$  is Poisson with rate  $s$ . By a union bound,

$$\mathbb{P} \left[ \bigcup_{u \in [0, s]} \mathcal{H}_i(t^* - u) \not\subseteq B(i, l) \right] \leq 2\mathbb{P}[\text{Po}(s) \geq l]. \quad (3.70)$$

The moment generating function of a poisson random variable with rate  $s$  is

$$M(t) = \exp(s(e^t - 1)). \quad (3.71)$$

Using a Chernoff bound we have for every  $t > 0$ ,

$$\mathbb{P}[\text{Po}(s) \geq l] \leq \exp(s(e^t - 1) - tl). \quad (3.72)$$

Overall we have

$$\mathbb{P} \left[ \bigcup_{u \in [0, s]} \mathcal{H}_i(t^* - u) \not\subseteq B(i, l) \right] \leq 2 \exp(s(e^t - 1) - tl). \quad (3.73)$$



Choosing  $t = \ln 2$ ,

$$\mathbb{P} \left[ \bigcup_{u \in [0, s]} \mathcal{H}_i(t^* - u) \not\subseteq B(i, l) \right] \leq 2 \exp(s - l \ln 2). \quad (3.74)$$

□

**Lemma 3.9.** *Let  $i$  and  $j$  be the indices of two vertices on the cycle  $(\mathbb{Z}/n\mathbb{Z})$  separated by distance  $k$ . Then*

$$\mathbb{P}[X_j = 1 | X_i = 1] \leq \frac{e^{-z}}{n} + 2 \left( \frac{2 - \sqrt{\theta(4 - \theta)}}{2 - \theta} \right)^k. \quad (3.75)$$

*Proof.* There are two ways in which the update history of vertex  $j$  can survive until time 0. The update history can survive without intersecting with the update history of vertex  $i$  or the update history of vertex  $j$  can merge with the update history of vertex  $i$  (whose survival we are conditioning on). Breaking up the probability this way we have

$$\begin{aligned} \mathbb{P}[X_j = 1 | X_i = 1] &= \mathbb{P}[X_j = 1, \mathcal{H}_i \cap \mathcal{H}_j = \emptyset | X_i = 1] \\ &\quad + \mathbb{P}[X_j = 1, \mathcal{H}_i \cap \mathcal{H}_j \neq \emptyset | X_i = 1] \end{aligned} \quad (3.76)$$

$$\leq \mathbb{P}[X_j = 1, \mathcal{H}_i \cap \mathcal{H}_j = \emptyset | X_i = 1] + \mathbb{P}[\mathcal{H}_i \cap \mathcal{H}_j \neq \emptyset | X_i = 1]. \quad (3.77)$$

The result follows from Lemmas 3.10 and 3.11. □

**Lemma 3.10.** *Let  $i$  and  $j$  be the indices of two vertices on the cycle  $(\mathbb{Z}/n\mathbb{Z})$ . Then*

$$\mathbb{P}[X_j = 1, \mathcal{H}_i \cap \mathcal{H}_j = \emptyset | X_i = 1] \leq e^{-z}/n. \quad (3.78)$$

*Proof.* To begin

$$\mathbb{P}[X_j = 1, \mathcal{H}_i \cap \mathcal{H}_j = \emptyset | X_i = 1] = \mathbb{P}[X_i = 1]^{-1} \mathbb{P}[X_i = 1, X_j = 1, \mathcal{H}_i \cap \mathcal{H}_j = \emptyset] \quad (3.79)$$

$$= \frac{\mathbb{P}[\mathcal{H}_i(0) \cup \mathcal{H}_j(0) \neq \emptyset, \mathcal{H}_i \cap \mathcal{H}_j = \emptyset]}{\mathbb{P}[\mathcal{H}_i(0) \neq \emptyset]}. \quad (3.80)$$

Define  $S$  to be the first time that the histories intersect, or define  $S = 0$  if the histories do not intersect before time 0. We construct a new history,  $\mathcal{H}'_j$ , in the following way. Along the interval  $[S, t^*)$ ,  $\mathcal{H}'_j$  is constructed from the update sequence in the exact same way as  $\mathcal{H}_j$ . However, along the interval  $[0, S)$ , we replace the update sequence with an another i.i.d. copy of the update sequence. We have constructed  $\mathcal{H}'_j$  to be independent of  $\mathcal{H}_i$ . This is obvious for  $t < S$ . For  $t \geq S$  we note that no updates are a part of both histories. From the memoryless property of the exponential waiting times between

updates, we get that each update is independent of the rest, and so the histories are independent too. Note that in the event  $\{\mathcal{H}_i \cap \mathcal{H}_j = \emptyset\}$ , we have  $\mathcal{H}_j = \mathcal{H}'_j$  and so

$$\mathbb{P}[\mathcal{H}_i(0) \cup \mathcal{H}_j(0) \neq \emptyset, \mathcal{H}_i \cap \mathcal{H}_j = \emptyset] = \mathbb{P}[\mathcal{H}_i(0) \cup \mathcal{H}'_j(0) \neq \emptyset, \mathcal{H}_i \cap \mathcal{H}_j = \emptyset] \quad (3.81)$$

$$\leq \mathbb{P}[\mathcal{H}_i(0) \cup \mathcal{H}'_j(0) \neq \emptyset] \quad (3.82)$$

$$= \mathbb{P}[\mathcal{H}_i(0) \neq \emptyset] \mathbb{P}[\mathcal{H}'_j(0) \neq \emptyset] \quad (3.83)$$

since  $\mathcal{H}_i$  and  $\mathcal{H}'_j$  are independent. Since  $\mathbb{P}[\mathcal{H}'_j(0) \neq \emptyset] = \mathbb{P}[\mathcal{H}_j(0) \neq \emptyset] = e^{-z}/n$  we get the desired result.  $\square$

**Lemma 3.11.** *Let  $i$  and  $j$  be the indices of two vertices that are separated by distance  $k$ . Then*

$$\mathbb{P}[\mathcal{H}_i \cap \mathcal{H}_j \neq \emptyset | X_j = 1] \leq 2 \left( \frac{2 - \sqrt{\theta(4 - \theta)}}{2 - \theta} \right)^k. \quad (3.84)$$

where  $k = |i - j|$ .

*Proof.* We first must deal with the effect that conditioning on  $X_j = 1$  has on the probability that the two update histories merge. For the history of vertex  $j$  to survive, all updates along the history must be non-oblivious updates. We also note that conditioning on the history of vertex  $j$  surviving should not result in an increase in the overall rate of updates (since each update is a chance that the history will die). By this reasoning,

$$\mathbb{P}[\mathcal{H}_i \cap \mathcal{H}_j \neq \emptyset | X_j = 1] \leq \mathbb{P}[\mathcal{H}_i \cap \bar{\mathcal{H}}_j] \quad (3.85)$$

where  $\bar{\mathcal{H}}_j$  is an undying random walk that starts at  $j$  and going backwards in time from  $t^*$  moves right at rate  $1/2$  and left at rate  $1/2$ .

We note that, while  $\mathcal{H}_i$  survives, the distance between  $\mathcal{H}_i$  and  $\bar{\mathcal{H}}_j$  is a birth and death process that starts at  $k = |i - j|$  and has birth and death rates,  $\lambda = \mu = (2 - \theta)/2$ . Let  $P(t)$  be such a process and define  $s_0 = \inf\{t : P(t) = 0\}$  to be the first time the process reaches zero (this corresponds to the update histories merging). Let  $s_d$  be exponentially distributed with rate  $\theta$  (this corresponds to the update history of vertex  $i$  dying). Then

$$\mathbb{P}[\mathcal{H}_i \cap \mathcal{H}_j \neq \emptyset | X_j = 1] \leq \mathbb{P}[\mathcal{H}_i \cap \bar{\mathcal{H}}_j] \leq 2\mathbb{P}_k(s_0 < s_d) \quad (3.86)$$

where  $\mathbb{P}_k$  indicates that  $P(0) = k$ . The factor of two comes from the fact that the update histories may meet by going the other direction around the cycle. We also note that we allow  $P(t)$  to continue beyond  $t = t^*$ , unlike our update histories which stop at time 0. This does not present a problem as the effect of allowing this is to increase the size of our upper bound.

At any time before  $s_d$  there are three possibilities for what can happen to  $P$  next. Either the next event is a birth with probability  $(2 - \theta)/4$ , the next event is a death with the same probability or we reach time  $s_d$  with probability  $\theta/2$ . Writing  $\zeta_k = \mathbb{P}_k(s_0 < s_d)$  this gives us the recurrence relation

$$\zeta_k = \frac{2 - \theta}{4}\zeta_{k-1} + \frac{2 - \theta}{4}\zeta_{k+1} \quad (3.87)$$

which is subject to the conditions

$$\zeta_0 = 1 \quad (3.88)$$

$$\zeta_k \leq 1, \forall k \in \mathbb{N}. \quad (3.89)$$

This recurrence has characteristic equation

$$x^2 - \frac{4}{2 - \theta}x + 1 = 0 \quad (3.90)$$

which has roots

$$r_1 = \frac{2 + \sqrt{\theta(4 - \theta)}}{2 - \theta} \quad (3.91)$$

$$r_2 = \frac{2 - \sqrt{\theta(4 - \theta)}}{2 - \theta} \quad (3.92)$$

and so

$$\zeta_k = ar_1^k + br_2^k \quad (3.93)$$

where  $a$  and  $b$  are constants to be determined from (3.88) and (3.89). We note that  $r_1 \geq 1, \forall \theta \in [0, 1]$  and so from (3.89) we have that  $a = 0$ . Finally from (3.88),  $b = 1$  and so

$$\zeta_k = \left( \frac{2 - \sqrt{\theta(4 - \theta)}}{2 - \theta} \right)^k. \quad (3.94)$$

□

## Chapter 4

# The Coupling Time on Vertex Transitive Graphs

For the most part, this chapter will be similar in structure and content to Chapter 3. The main difference is that we extend the family of graphs on which we consider the Ising heat-bath Glauber dynamics from the cycle to any vertex transitive graph. Again, the main result, Theorem 4.1, concerns the coupling time,  $T_n$ , as defined in Section 2.2.2, and establishes that at sufficiently high temperature (that is, for  $\beta$  small enough), the coupling time converges in distribution to a Gumbel distribution.

Restricting  $\beta$  to be sufficiently small is a consequence of the increased generality of this chapter. As mentioned in Chapter 3, in the high-temperature regime we expect the dynamics to be similar to those when  $\beta = 0$ . When  $\beta = 0$  the problem simplifies to the coupon collector's problem, which is known to have a Gumbel limit. However, at the critical temperature, and below in the low-temperature regime, there is no reason to suspect that the dynamics will behave similarly to when  $\beta = 0$ . So our restriction of  $\beta$  to be small enough for the result to hold is, on at least a descriptive level, somewhat expected.

Our result partially confirms the conjecture by Collevocchio et al. in [17] that the coupling time for the Ising heat-bath process on the  $d$ -dimensional lattice,  $G_L = (\mathbb{Z}/L\mathbb{Z})^d$ , converges to a Gumbel distribution as  $L \rightarrow \infty$  for all  $\beta < \beta_C$ . Our result does not hold all the way up until the critical temperature. This is due to the fact that we are considering a larger class of graphs than just the square lattice, and so it is unreasonable to expect such a result to provide sharp bounds for the lattice. A separate treatment of the square lattice in particular may be needed for a result holding all the way up to the critical temperature. Note that this is what Lubetzky and Sly did in [14] to prove the existence of cutoff for the full high-temperature regime. Since our proof is also based on

information percolation there is good reason to think that a similar approach could also work to extend our result.

[MENTION ONE LEMMA WHERE WE NEED BETA SMALL IS VERY SIMILAR TO THEIR LEMMA?]

To state the main result we first must define the graphs on which it is valid. Let  $(G_n)$  be a sequence of vertex-transitive graphs with  $n$  vertices. Let  $P_n(k)$  denote the number of vertices at distance  $k$  from a vertex  $i$  in  $G_n$  and let  $Q_n(k)$  denote the number of vertices at distance  $k$  or less from a vertex  $i$  in  $G_n$ . Define

$$\mathcal{G} = \left\{ (G_n) : \forall n \Delta_n = C_1, Q_n(\ln^2(n))/n \rightarrow 0, \lim_{n \rightarrow \infty} \sum_{k=1}^{\infty} P_n(k) e^{-k} = C_2, C_1, C_2 > 0 \right\} \quad (4.1)$$

where  $\Delta_n$  is the vertex degree of  $G_n$ . In this chapter we consider sequences of vertex-transitive graphs  $(G_n) \in \mathcal{G}$ . Note that  $P$  and  $Q$  polynomial in  $k$  is enough for  $(G_n)$  to be in  $\mathcal{G}$ . This includes the  $d$ -dimensional square lattices.

[MAKE THIS BETTER AFTER TIM HAS GOT BACK ABOUT THIS]

We now define a couple of quantities. Firstly, the *magnetization at vertex  $i$  at time  $t$*  is

$$m_t(i) = \mathbb{E}[\mathcal{T}_t[i]] \quad (4.2)$$

where  $(\mathcal{T}_t)_{t \geq 0}$  is the dynamics starting from the all-plus configuration. Note that on transitive graphs, with which this chapter is concerned, we can drop the dependence on  $i$  and simply write  $m_t$  for the magnetization at any vertex at time  $t$ . This quantity is not to be confused with the magnetization on a volume (as mentioned in Section 2.1.1) which has nothing to do with dynamics. The magnetization on a volume is a random variable given by the sum of all the spins on a given volume where the spins are distributed according to the Gibbs measure. In contrast, the magnetization at a vertex, as defined here, is a deterministic function of time that is a property of the Glauber dynamics.

We can now define the time

$$t_c(n) = \inf \left\{ t > 0 : m_t = \frac{1}{n} \right\}. \quad (4.3)$$

which is around the time it takes for the top and bottom chains to couple.

**Theorem 4.1.** *Let  $(G_n) \in \mathcal{G}$  be a sequence of vertex-transitive graphs,  $G_n = (V, E)$  with  $|V| = n$  vertices. Let  $T_n$  be the coupling time for the continuous-time Ising heat-bath dynamics for the zero-field ferromagnetic Ising model on  $G_n$ . Then for any small*

enough inverse-temperature  $\beta$ ,

$$\lim_{n \rightarrow \infty} \mathbb{P}[T_n < z + t_c(n)] = e^{-e^{-z}}. \quad (4.4)$$

The proof of Theorem 4.1 will be given in Section 4.3 after the essential preliminaries are presented. In Section 4.1 we describe an alternative construction of the histories that can sometimes be easier to work with. Then in Section 4.2, we outline the overall approach to the proof. The method is very similar to the method used in Chapter 3 but there are some additional problems that are addressed. Finally, we defer results directly concerning the update histories to Section 4.4.

## 4.1 Information Percolation in higher dimensions

In the previous chapter, we showed that on the cycle, there was a coupling that made the update history of a single vertex to be a continuous-time random walk that died at rate  $\theta$ . On lattices of dimension  $d > 2$ , we can no longer use this coupling and so the updates histories are significantly more complex.

Recall from Section 2.3.2.2 that given a target time  $t^*$ , the update history of a vertex set  $A$  at time  $t$ ,  $\mathcal{H}_A(t)$ , is the set of vertices whose spins at time  $t$  determine the spins of  $A$  at time  $t^*$ . Developing this history backwards in time from  $t = t^*$  produces a subgraph of  $\Omega \times [0, t^*]$  which we write as  $\mathcal{H}_A$  and call the update history of vertex set  $A$ . This history can be constructed using the update sequence along  $(t, t^*]$ .

In practise, we may choose to construct this history as follows: For each  $i \in A$ , create a temporal edge between  $(i, t^*)$  and  $(i, t_i)$  where  $t_i$  is the time of the latest update to  $i$  (or 0 if  $i$  is never updated). Then for each update  $(i, u, t_i)$ , we either terminate the edge if  $u$  is such that the update is oblivious, or we add spatial branches to each of the neighbours of  $i$ . We repeat this process recursively for the neighbours of  $i$  until every branch has been terminated due to an oblivious update or has reached time 0.

However, it is possible for vertices to be removed from  $\mathcal{H}_A(t)$  from updates that are not oblivious (see Figure 2.4). Since our method above for constructing the history does not take this into account, the history it produces will possibly be larger than  $\mathcal{H}_A$ . To ensure a distinction between the two, the history that results from the above construction we will denote  $\hat{\mathcal{H}}_A$ , and likewise  $\hat{\mathcal{H}}_A(t)$  for the history at time  $t$  that results from the above construction. We have that

$$\mathcal{H}_A(t) \subseteq \hat{\mathcal{H}}_A(t) \quad (4.5)$$

and also that  $\mathcal{H}_A$  is a subgraph of  $\hat{\mathcal{H}}_A$ .

### 4.1.1 Magnetization

One quantity which we used multiple times in Chapter 3 was  $\mathbb{P}[X_i = 1]$ . Although it was not required earlier, we would now like to make clear that this is in fact the magnetization at time  $t^*$ .

Recall that the magnetization at vertex  $i \in V$  at time  $t > 0$  is defined to be

$$m_t(i) = \mathbb{E}[\mathcal{T}_t[i]] \quad (4.6)$$

where  $(\mathcal{T}_t)_{t \geq 0}$  is the dynamics starting from the all-plus configuration. Given a monotonically coupled chain  $(\mathcal{B}_t)_{t \geq 0}$ , starting in the all minus configuration and such that  $\mathcal{T}_t[i] \geq \mathcal{B}_t[i]$  for all  $t \geq 0$  and  $i \in V$ , we can split up this expectation by conditioning on the event  $A_t = \{\mathcal{T}_t[i] \neq \mathcal{B}_t[i]\}$ . We obtain that

$$m_t(i) = \mathbb{E}[Y_t^+[i]] \quad (4.7)$$

$$\begin{aligned} &= \mathbb{P}[A_t] (\mathbb{P}[Y_t^+[i] = 1|A_t] - \mathbb{P}[Y_t^+[i] = -1|A_t]) \\ &\quad + \mathbb{P}[A_t^c] (\mathbb{P}[Y_t^+[i] = 1|A_t^c] - \mathbb{P}[Y_t^+[i] = -1|A_t^c]). \end{aligned} \quad (4.8)$$

Now if event  $A_t^c$  holds,  $\mathcal{T}_t[i] = \mathcal{B}_t[i]$ , and so by symmetry vertex  $i$  must take values  $-1$  and  $+1$  uniformly. Furthermore, by the monotonicity of our coupling, if  $A_t$  holds, we must have that  $\mathcal{T}_t[i] = +1$  and  $\mathcal{B}_t[i] = -1$ . So

$$m_t(i) = \mathbb{P}[A_t]. \quad (4.9)$$

Finally, given a target time  $t^*$ ,  $X_i$  is defined such that  $\{X_i = 1\} = A_{t^*}$ . So

$$\mathbb{P}[X_i = 1] = m_{t^*}(i). \quad (4.10)$$

We end this section with some results concerning the magnetization, and in particular, the magnetization at time

$$t^* = t_c(n) + z. \quad (4.11)$$

The following comes from [15] and is valid on any graph, not just transitive ones.

**Lemma 4.2** ([15], Claim 3.3). *On any graph with maximum degree  $\Delta$ , for any  $t, s > 0$  we have*

$$e^{-2s} \leq \frac{\sum_i m_{t+s}[i]^2}{\sum_i m_t[i]^2} \leq e^{-2(1-\beta\Delta)s}. \quad (4.12)$$

The following corollaries are then straightforward.

**Corollary 4.3.** *On any vertex transitive graph with degree  $\Delta$ , for any  $t, s > 0$  we have*

$$e^{-s}m_t \leq m_{t+s} \leq m_t e^{-(1-\beta\Delta)s}. \quad (4.13)$$

**Corollary 4.4.** *On any vertex transitive graph with degree  $\Delta$ ,  $m_{t^*}$  can be bounded as follows:*

For  $z \geq 0$ ,

$$\frac{e^{-z}}{n} \leq m_{t^*} \leq \frac{e^{-(1-\beta\Delta)z}}{n}. \quad (4.14)$$

For  $z \leq 0$ ,

$$\frac{e^{-(1-\beta\Delta)z}}{n} \leq m_{t^*} \leq \frac{e^{-z}}{n}. \quad (4.15)$$

Bearing in mind that  $m_0 = 1$ , we also obtain a bound on  $t_c(n)$ .

**Corollary 4.5.** *On any vertex transitive graph with degree  $\Delta$ , for  $\beta < 1/\Delta$*

$$\ln(n) \leq t_c(n) \leq \frac{\ln(n)}{1 - \beta\Delta}. \quad (4.16)$$

## 4.2 Setup

As in Chapter 3, we prove Theorem 4.1 through the use of a stronger statement in Theorem 4.6. The overall approach is almost exactly as described in Section 3.2; at a time  $t^*$  we count the number of vertices at which the bottom and top chains differ and show that the distribution of this random variable, which we call  $W$ , is close to an appropriately chosen compound poisson distribution. The main difference in the construction is that we use a different  $t^*$  here.

Recalling the definition of  $t_c(n)$  in (4.3), fix  $z$  and a time of interest  $t^* = t_c(n) + z$ . From here we define  $X_i$ ,  $U_i$ ,  $W_i$ ,  $\delta_1$ ,  $\delta_4$ ,  $\lambda$ , and  $\mu$  exactly as in Chapter 3 but using our new definition for  $t^*$ . For convenience, we repeat the definitions here.

For each vertex  $i \in V$ , define indicators

$$X_i = \begin{cases} 1 & \mathcal{B}_{t^*}[i] \neq \mathcal{T}_{t^*}[i], \\ 0 & \mathcal{B}_{t^*}[i] = \mathcal{T}_{t^*}[i] \end{cases} \quad (4.17)$$



and set  $W = \sum_{i \in V} X_i$ . For each  $i \in V$ , decompose  $W$  into  $W = X_i + U_i + Z_i + W_i$  where

$$U_i = \sum_{j \in B_i} X_j, \quad Z_i = \sum_{j \in C_i} X_j, \quad W_i = \sum_{j \in D_i} X_j. \quad (4.18)$$

and  $B_i, C_i$ , and  $D_i$  are the vertex sets

$$B_i = \{j \neq i : |j - i| \leq b_n\}, \quad (4.19)$$

$$C_i = \{j \notin B_i \cup \{i\} : |j - i| \leq c_n\}, \quad (4.20)$$

$$D_i = V \setminus (B_i \cup C_i \cup \{i\}). \quad (4.21)$$

As in Chapter 3, we have some freedom in choosing  $b_n$  and  $c_n$  but we will again choose  $b_n = \ln(n)$  and  $c_n = \ln(n)^2$ .

We now define the quantities

$$\lambda = \sum_{i \in V} \mathbb{E} \left[ \frac{X_i}{X_i + U_i} I[X_i + U_i \geq 1] \right], \quad (4.22)$$

$$\mu_l = \frac{1}{l\lambda} \sum_{i \in V} \mathbb{E} [X_i I[X_i + U_i = l]], \quad l \geq 1, \quad (4.23)$$

which will be the parameters of the approximating compound Poisson distribution to  $W$ . We also define

$$\delta_1 = \sum_{i \in V} \sum_{k \geq 0} \mathbb{P}[X_i = 1, U_i = k] \mathbb{E} \left| \frac{\mathbb{P}[X_i = 1, U_i = k | W_i]}{\mathbb{P}[X_i = 1, U_i = k]} - 1 \right|, \quad (4.24)$$

$$\delta_4 = \sum_{i \in V} (\mathbb{E}[X_i Z_i] + \mathbb{E}[X_i] \mathbb{E}[X_i + U_i + Z_i]), \quad (4.25)$$

which we desire to be small for the compound Poisson approximation to be good.

The following theorem (reworked from [19]) bounds the distance between the distributions of  $W$  and the approximating compound Poisson.

**Theorem 4.6** ([19]). *Let  $W$ ,  $\lambda$ ,  $\mu$ ,  $\delta_1$  and  $\delta_4$  be as defined above. Then*

$$d_{\text{TV}}(\mathcal{L}(W), \text{CP}(\lambda, \mu)) \leq (\delta_1 + \delta_4) e^\lambda. \quad (4.26)$$

As per the discussion proceeding Theorem 3.2, we obtain the following as a corollary to Theorem 4.6.

**Corollary 4.7.** *Let  $T_n$  be the coupling time of the continuous-time heat-bath Glauber dynamics for the zero-field Ising model at inverse-temperature  $\beta$  on the graph  $G_n$  and*

let  $\delta_1$ ,  $\delta_4$  and  $\lambda$  be as defined above. Then

$$\left| \mathbb{P}[T_n \leq z + t_c(n)] - e^{-\lambda} \right| \leq (\delta_1 + \delta_4)e^\lambda, \quad (4.27)$$

where  $\theta = 1 - \tanh(2\beta)$ .

### 4.3 Proof of Theorem 4.1

In this section we use Corollary 4.7 to prove Theorem 4.1 by bounding  $\lambda$  and showing that  $\delta_1$  and  $\delta_4$  go to zero as  $n \rightarrow \infty$ . This is done in Lemmas 4.8, 4.9, and 4.10. The proofs of these require some additional lemmas concerning properties of the update histories which have been deferred to Section 4.4.

**Lemma 4.8.** *Using the above setup, there exists a constant  $C \in (0, 1)$  such that*

$$C \min(e^{-z}, e^{-(1-\beta\Delta)z}) \leq \lim_{n \rightarrow \infty} \lambda \leq \max(e^{-z}, e^{-(1-\beta\Delta)z}). \quad (4.28)$$

*Proof.* Starting with the definition of  $\lambda$ , we have

$$\lambda = \sum_{i \in V} \mathbb{E} \left[ \frac{X_i}{X_i + U_i} I[X_i + U_i \geq 1] \right] \quad (4.29)$$

$$= \sum_{i=1}^n \mathbb{P}(X_i = 1) \mathbb{E} \left[ \frac{1}{1 + U_i} | X_i = 1 \right] \quad (4.30)$$

$$= nm_{t^*} \mathbb{E} \left[ \frac{1}{1 + U_i} | X_i = 1 \right] \quad (4.31)$$

where we have used that  $X_i$  is zero-one, (4.10), and the transitivity of the graph. Clearly

$$\mathbb{E} \left[ \frac{1}{1 + U_i} | X_i = 1 \right] \leq 1 \quad (4.32)$$

and so  $\lambda \leq nm_{t^*} \leq \max(e^{-z}, e^{-(1-\beta\Delta)z})$ .

By Jensen's inequality

$$\mathbb{E} \left[ \frac{1}{1 + U_i} | X_i = 1 \right] \geq \frac{1}{\mathbb{E}[1 + U_i | X_i = 1]} \quad (4.33)$$

$$= \frac{1}{1 + \mathbb{E}[U_i | X_i = 1]}. \quad (4.34)$$

so in order to find a lower bound for  $\lambda$  we will find an upper bound to  $\mathbb{E}[U_i|X_i = 1]$ . Now by Lemma 4.14, there exists a  $C_1 > 0$  such that for small enough  $\beta$ ,

$$\mathbb{E}[U_i|X_i = 1] = \sum_{j \in B_i} \mathbb{P}[X_j = 1|X_i = 1] \quad (4.35)$$

$$\leq |B_i|m_{t^*} + C_1 \sum_{k=1}^{\lfloor b_n \rfloor} \sum_{|j-i|=k} e^{-k}. \quad (4.36)$$

From (4.1),

$$\mathbb{E}[U_i|X_i = 1] \leq |B_i|m_{t^*} + C_1 \sum_{k=1}^{\lfloor b_n \rfloor} P_n(k)e^{-k} \quad (4.37)$$

$$\leq |B_i|m_{t^*} + C_1 \sum_{k=1}^{\infty} P_n(k)e^{-k} \quad (4.38)$$

$$= C_z|B_i|/n + C_2 \quad (4.39)$$

for some  $C_2 > 0$ ,  $C_z = \max(e^{-z}, e^{-(1-\beta\Delta)z})$ . As  $n \rightarrow \infty$ , the first term vanishes and we are left with

$$\lim_{n \rightarrow \infty} \lambda \geq \frac{1}{1 + C_2} nm_{t^*} \quad (4.40)$$

$$\geq C \min(e^{-z}, e^{-(1-\beta\Delta)z}) \quad (4.41)$$

for some  $C \in (0, 1)$ . □

**Lemma 4.9.** *Let  $\delta_1$  be as defined above in (4.24). Then*

$$\lim_{n \rightarrow \infty} \delta_1 = 0. \quad (4.42)$$

*Proof.* Starting with the definition of  $\delta_1$ , we have

$$\delta_1 = \sum_{i=1}^n \sum_{k=0}^{|B_i|} \mathbb{P}[X_i = 1, U_i = k] \mathbb{E} \left| \frac{\mathbb{P}[X_i = 1, U_i = k|W_i]}{\mathbb{P}[X_i = 1, U_i = k]} - 1 \right| \quad (4.43)$$

$$= n \sum_{k=0}^{|B_i|} \mathbb{E} |\mathbb{P}[X_i = 1, U_i = k|W_i] - \mathbb{P}[X_i = 1, U_i = k]| \quad (4.44)$$

by the transitivity of the graph. Let

$$C_i^c = \{j : |j - i| \leq (c_n + b_n)/2\} \quad (4.45)$$

be the set of vertices within distance  $(b_n + c_n)/2$  of  $i$  and define the events

$$A_1 = \{\exists j \in B_i \cup \{i\}, \exists t \in [0, t^*] : \mathcal{H}_j(t) \not\subseteq C_i^c\} \quad (4.46)$$

and

$$A_2 = \{\exists j \in D_i, \exists t \in [0, t^*] : \mathcal{H}_j(t) \cap C_i^c \neq \emptyset\} \quad (4.47)$$

as well as their intersection

$$A = A_1 \cap A_2. \quad (4.48)$$

From Lemma 3.7,

$$\mathbb{P}[X_i = 1, U_i = j | A^c, W_i] = \mathbb{P}[X_i = 1, U_i = j | A^c]. \quad (4.49)$$

Continuing on from (4.44), we split the probabilities into

$$\delta_1 = n \sum_{k=0}^{|B_i|} \mathbb{E} \left| \mathbb{P}[X_i = 1, U_i = k | W_i, A] \mathbb{P}[A | W_i] - \mathbb{P}[X_i = 1, U_i = k | A] \mathbb{P}[A] + \right. \quad (4.50)$$

$$\left. \mathbb{P}(X_i = 1, U_i = k | A^c) (\mathbb{P}[A^c | W_i] - \mathbb{P}[A^c]) \right|$$

$$\leq n(|B_i| + 1) \mathbb{E} \left[ \mathbb{P}[A | W_i] + \mathbb{P}[A] + \left| \mathbb{P}[A^c | W_i] - \mathbb{P}[A^c] \right| \right] \quad (4.51)$$

$$= n(|B_i| + 1) \mathbb{E} [\mathbb{P}[A | W_i] + \mathbb{P}[A] + |1 - \mathbb{P}[A | W_i] - (1 - \mathbb{P}[A])|] \quad (4.52)$$

$$\leq n(|B_i| + 1) \mathbb{E} [\mathbb{P}[A | W_i] + \mathbb{P}[A] + \mathbb{P}[A | W_i] + \mathbb{P}[A]] \quad (4.53)$$

$$= 2n(|B_i| + 1) (\mathbb{E}[\mathbb{P}[A | W_i]] + \mathbb{P}[A]) \quad (4.54)$$

$$= 4n(|B_i| + 1) \mathbb{P}[A] \quad (4.55)$$

For either  $A_1$  or  $A_2$  to hold, there must exist a history that spreads at least distance  $(c_n - b_n)/2$  away from its starting vertex. By a union bound

$$\mathbb{P}[A] \leq \sum_{j=1}^n \mathbb{P}[\mathcal{H}_i \not\subseteq B(i, (c_n - b_n)/2) \times [0, t^*]] \quad (4.56)$$

$$= n \mathbb{P} \left[ \bigcup_{u \in [0, t^*]} \mathcal{H}_i(t^* - u) \not\subseteq B(i, (c_n - b_n)/2) \right] \quad (4.57)$$

Combining this with Lemma 4.12, and recalling our choices of  $b_n = \ln(n)$  and  $c_n = \ln(n)^2$  we get that

$$\delta_1 \leq 4n^2(|B_i| + 1) \exp(t^* \Delta - \ln \Delta (c_n - b_n)/2) \quad (4.58)$$

$$\leq 4n^{2+\Delta/(1-\beta\Delta)} (|B_i| + 1) \exp(\Delta z) \exp(-\ln \Delta (c_n - b_n)/2) \quad (4.59)$$

which goes to 0 as  $n \rightarrow \infty$ .

□

**Lemma 4.10.** *Let  $\delta_4$  be as defined above in (4.25). Then*

$$\lim_{n \rightarrow \infty} \delta_4 = 0. \quad (4.60)$$

*Proof.* Starting with the definition of  $\delta_4$ , we have

$$\delta_4 = \sum_{i=1}^n (\mathbb{E}[X_i Z_i] + \mathbb{E}[X_i] \mathbb{E}[X_i + U_i + Z_i]) \quad (4.61)$$

$$= n \mathbb{E}[X_i Z_i] + n m_{t^*}^2 (1 + |B_i| + |C_i|) \quad (4.62)$$

$$= n m_{t^*} \mathbb{E}[Z_i | X_i = 1] + n m_{t^*}^2 (1 + |B_i| + |C_i|) \quad (4.63)$$

$$\leq C_z \mathbb{E}[Z_i | X_i = 1] + \frac{C_z^2}{n} (1 + |B_i| + |C_i|) \quad (4.64)$$

where

$$C_z = \max(e^{-z}, e^{-(1-\beta\Delta)z}). \quad (4.65)$$

Now from (4.1), the second term above vanishes as  $n \rightarrow \infty$ . So we turn our attention to the first term. From Lemma 4.14, there exists a  $C > 0$  such that for small enough  $\beta$ ,

$$\mathbb{E}[Z_i | X_i = 1] = \sum_{j \in C_i} \mathbb{P}[X_j = 1 | X_i = 1] \quad (4.66)$$

$$\leq |C_i| (m_{t^*} + C e^{-b_n}) \quad (4.67)$$

$$\leq |C_i| \left( \frac{C_z}{n} + \frac{C}{n} \right) \quad (4.68)$$

which goes to zero as  $n \rightarrow \infty$ .

□

## 4.4 Additional Lemmas

The first of our additional Lemmas comes from [21, Lemma 3.1]. We have made our statement of the lemma slightly more precise and so we have rewritten both the lemma and proof out here along with our modifications. In particular, we have specified precisely how small  $\beta$  must be for the statement to hold.

The Lemma concerns two quantities,  $\chi(\mathcal{H}_A)$  and  $\mathcal{L}(\mathcal{H}_A)$  which in some sense measure the horizontal and vertical size of  $\mathcal{H}_A$  respectively. Define

$$\chi(\mathcal{H}_i) = \# \{((u, t), (v, t)) \in \mathcal{H}_i\} \quad (4.69)$$

which counts the total number of spatial edges in  $\mathcal{H}_i$  and define

$$\mathcal{L}(\mathcal{H}_i) = \sum_{i \in V} \int_0^{t^*} I_{(i,t) \in \mathcal{H}_i} dt \quad (4.70)$$

which is the sum of the lengths of all the temporal edges in  $\mathcal{H}_i$ .

**Lemma 4.11** ([21]). *For any  $0 \leq \eta < 1$ ,  $\lambda \in \mathbb{R}$ ,  $\alpha > -\ln(1 - \eta)$ , if*

$$\tanh(\beta\Delta) \leq \frac{1 - \eta - e^{-\alpha}}{e^{(\alpha+\lambda)\Delta} - e^{-\alpha}}, \quad (4.71)$$

*then for any  $A \subseteq V$ ,*

$$\mathbb{E}[\exp(\lambda\chi(\mathcal{H}_A) + \eta\mathcal{L}(\mathcal{H}_A))] \leq \exp(\alpha|A|) \quad (4.72)$$

*where  $\mathcal{H}_A$  is defined as*

$$\mathcal{H}_A = \bigcup_{i \in A} \mathcal{H}_i. \quad (4.73)$$

*Proof.* We first relax our histories to our alternative construction by observing that

$$\chi(\mathcal{H}_i) \leq \chi(\hat{\mathcal{H}}_i), \quad Z(\mathcal{H}_i) \leq Z(\hat{\mathcal{H}}_i). \quad (4.74)$$

Let  $W_s = |\hat{\mathcal{H}}_i(t^* - s)|$ , let  $Y_s = \chi(\hat{\mathcal{H}}_i \cap V \times [t^* - s, t^*])$  count the total number of spatial edges observed in the history by time  $t^* - s$  and let  $Z_s = \mathcal{L}(\hat{\mathcal{H}}_i \cap V \times [t^* - s, t^*])$ .

Initially,  $W_0 = 1$ ,  $Y_0 = 0$ , and  $Z_0 = 0$ . Recall that an oblivious update of a vertex causes it to be removed from the history and that a non-oblivious update causes the history to branch out to its  $\Delta$  neighbours. Oblivious updates occur at rate  $\theta W_s$  and cause  $W_s$  to decrease by 1. Non-oblivious updates occur at rate  $(1 - \theta)W_s$  and cause both  $W_s$  and  $Y_s$  to increase by no more than  $\Delta$ . The length,  $Z_s$ , grows as  $dZ_s = W_s ds$ . Therefore we can create a coupled process  $(\bar{W}_s, \bar{Y}_s, \bar{Z}_s)$  such that  $\bar{W}_s \geq W_s$ ,  $\bar{Y}_s \geq Y_s$ , and  $\bar{Z}_s \geq Z_s$  in the following way. We start with  $(\bar{W}_s, \bar{Y}_s, \bar{Z}_s) = (|A|, 0, 0)$  and at rate  $\theta\bar{W}_s$ ,  $\bar{W}_s$  decreases by 1; at rate  $(1 - \theta)\bar{W}_s$ , both  $\bar{W}_s$  and  $\bar{Y}_s$  increase by  $\Delta$ ; and  $\bar{Z}_s$  grows as  $d\bar{Z}_s = \bar{W}_s ds$ .

Let  $Q_s = \exp(\alpha\bar{W}_s + \lambda\bar{Y}_s + \eta\bar{Z}_s)$  where  $\alpha$ ,  $\lambda$ , and  $\eta$  are some fixed constants yet to be determined, and  $\alpha > -\ln(1 - \eta)$ . We will show that  $Q_s$  is a supermartingale. Let  $h$  be

some small time-step. Then [FIX THIS]

$$\begin{aligned}
\mathbb{E}[Q_{s_0+h} - Q_{s_0} | Q_{s_0}] &= h\theta\bar{W}_{s_0} (\exp(\alpha(\bar{W}_{s_0} - 1) + \lambda\bar{Y}_{s_0}) - \exp(\alpha\bar{W}_{s_0} + \lambda\bar{Y}_{s_0})) + \quad (4.75) \\
&\quad h(1-\theta)\bar{W}_{s_0} (\exp(\alpha(\bar{W}_{s_0} + \Delta) + \lambda(\bar{Y}_{s_0} + \Delta)) - \exp(\alpha\bar{W}_{s_0} + \lambda\bar{Y}_{s_0})) + \\
&\quad \mathcal{O}(h^2) \\
&= \left( \eta + \theta(e^{-\alpha} - 1) + (1-\theta)(e^{(\alpha+\lambda)\Delta} - 1) \right) h\bar{W}_{s_0}Q_{s_0} + \mathcal{O}(h^2). \quad (4.76)
\end{aligned}$$

Dividing through by  $h$  and taking  $h$  to 0, we have

$$\left. \frac{d}{ds} \mathbb{E}[Q_s | Q_{s_0}] \right|_{s=s_0} = \left( \eta + \theta(e^{-\alpha} - 1) + (1-\theta)(e^{(\alpha+\lambda)\Delta} - 1) \right) \bar{W}_{s_0}Q_{s_0} \quad (4.77)$$

which is non-positive when

$$\theta \geq \frac{\eta + e^{(\alpha+\lambda)\Delta} - 1}{e^{(\alpha+\lambda)\Delta} - e^{-\alpha}} \quad (4.78)$$

or in terms of the inverse temperature,  $\beta$ ,

$$\tanh(\beta\Delta) \leq \frac{1 - \eta - e^{-\alpha}}{e^{(\alpha+\lambda)\Delta} - e^{-\alpha}}. \quad (4.79)$$

Hence  $Q_s$  is a supermartingale when (4.79) holds. Define the stopping time

$$\tau = \inf\{s : \bar{W}_s = 0\}. \quad (4.80)$$

At this time, the histories have completely died out. Finally, from optional stopping [FIGURE THIS OUT PROPERLY],

$$\mathbb{E}[\exp(\lambda\bar{Y}_\tau + \eta\bar{Z}_\tau)] \leq \mathbb{E}[Q_0] \quad (4.81)$$

$$= \exp(\alpha|A|). \quad (4.82)$$

□

This next Lemma is based on [11, Lemma 2.1] and bounds the speed at which the histories can spread through the graph.

**Lemma 4.12.** *Let  $B(i, l)$  indicate the set of vertices at distance  $l$  or smaller from vertex  $i$ . The probability that the history of vertex  $i$  escapes  $B(i, l)$  in time  $s$  is bounded by*

$$\mathbb{P} \left[ \bigcup_{u \in [0, s]} \mathcal{H}_i(t^* - u) \not\subseteq B(i, l) \right] \leq \exp(s\Delta^2 - l \ln \Delta). \quad (4.83)$$

*Proof.* Let  $\mathcal{W} = \{w = (w_1, w_2, \dots, w_l) : w_1 = i, ||w_{k-1} - w_k|| = 1\}$  be the set of length  $l$  sequences of adjacent vertices starting at vertex  $i$ . If  $\mathcal{H}_i$  contains any vertex outside  $B(i, l)$  at a time  $u \in [t^* - s, t^*]$  then there must be some sequence  $w \in \mathcal{W}$  such that each  $w_i$  was updated at some time  $t^* > t_i > t^* - s$  and  $t_{k-1} > t_k$ . Call this event  $M_w$ . For any particular sequence  $w$ ,

$$\mathbb{P}[M_w] = \mathbb{P}[\text{Po}(s) \geq l] \quad (4.84)$$

where  $\text{Po}(s)$  is Poisson with rate  $s$ . By a union bound over  $\mathcal{W}$ ,

$$\mathbb{P} \left[ \bigcup_{u \in [0, s]} \mathcal{H}_i(t^* - u) \not\subseteq B(i, l) \right] \leq \Delta^{l-1} \mathbb{P}[\text{Po}(s) \geq l]. \quad (4.85)$$

The moment generating function of a poisson random variable with rate  $s$  is

$$M(t) = \exp(s(e^t - 1)). \quad (4.86)$$

Using a Chernoff bound we have for every  $t > 0$ ,

$$\mathbb{P}[\text{Po}(s) \geq l] \leq \exp(s(e^t - 1) - tl). \quad (4.87)$$

Overall we have

$$\mathbb{P} \left[ \bigcup_{u \in [0, s]} \mathcal{H}_i(t^* - u) \not\subseteq B(i, l) \right] \leq \Delta^{l-1} \exp(s(e^t - 1) - tl) \quad (4.88)$$

$$\leq \exp(s(e^t - 1) + l(\ln \Delta - t)). \quad (4.89)$$

Choosing  $t = 2 \ln \Delta$ ,

$$\mathbb{P} \left[ \bigcup_{u \in [0, s]} \mathcal{H}_i(t^* - u) \not\subseteq B(i, l) \right] \leq \exp(s(\Delta^2 - 1) - l \ln \Delta) \quad (4.90)$$

$$\leq \exp(s\Delta^2 - l \ln \Delta). \quad (4.91)$$

□

[IF WE CAN SOMEHOW SHOW THAT  $\mathbb{P}[\mathcal{H}_i \cap \mathcal{H}_j \neq \emptyset | X_1 = 1] \leq \mathbb{P}[\tilde{\mathcal{H}}_i \cap \mathcal{H}_j]$  THEN THIS GIVES US A BIGGER RANGE OF TEMPERATURES.]

**Lemma 4.13.** *blah*



*Proof.* Write  $\bar{\mathcal{H}}_i$  for the history at  $i$  but it only has branching updates at rate 1. Write  $d_j$  for the time of death of history  $j$ . Let  $k = |i - j|$ .

$$\begin{aligned} \mathbb{P}(\bar{\mathcal{H}}_i \cap \mathcal{H}_j) &= \mathbb{P}(\bar{\mathcal{H}}_i \cap \mathcal{H}_j | d_j < c(k)) \mathbb{P}(d_j < c(k)) \\ &\quad + \mathbb{P}(\bar{\mathcal{H}}_i \cap \mathcal{H}_j | d_j \geq c(k)) \mathbb{P}(d_j \geq c(k)) \end{aligned} \quad (4.92)$$

$$\leq \mathbb{P}(\bar{\mathcal{H}}_i \cap \mathcal{H}_j | d_j < c(k)) + \mathbb{P}(d_j \geq c(k)) \quad (4.93)$$

$$\leq \mathbb{P}(\bar{\mathcal{H}}_i \cap \mathcal{H}_j | d_j < c(k)) + e^{-(1-\beta\Delta)c(k)} \quad (4.94)$$

$$\leq 2\mathbb{P}(\bar{\mathcal{H}}_i(t^* - c(k)) \not\subseteq B(i, k/2)) + e^{-(1-\beta\Delta)c(k)} \quad (4.95)$$

$$\leq \exp(c(k)\Delta^2 - k/2 \ln \Delta) + e^{-(1-\beta\Delta)c(k)} \quad (4.96)$$

$$(4.97)$$

Choosing  $c(k) = \ln(k)^2$  works so long as  $\beta\Delta < 1$ .  $\square$

**Lemma 4.14.** *There exists a constant  $C > 0$  such that for sufficiently small  $\beta$ ,*

$$\mathbb{P}[X_j = 1 | X_i = 1] \leq m_{t^*} + Ce^{-k}. \quad (4.98)$$

where  $k = |i - j|$  is the distance between vertices  $i$  and  $j$ .

*Proof.* There are two ways in which the update history of vertex  $j$  can survive until time 0. The update history can survive without intersecting with the update history of vertex  $i$  or the update history of vertex  $j$  can merge with the update history of vertex  $i$  (whose survival we are conditioning on). Breaking up the probability this way we have

$$\begin{aligned} \mathbb{P}[X_j = 1 | X_i = 1] &= \mathbb{P}[X_j = 1, \mathcal{H}_i \cap \mathcal{H}_j = \emptyset | X_i = 1] \\ &\quad + \mathbb{P}[X_j = 1, \mathcal{H}_i \cap \mathcal{H}_j \neq \emptyset | X_i = 1] \end{aligned} \quad (4.99)$$

$$\leq \mathbb{P}[X_j = 1, \mathcal{H}_i \cap \mathcal{H}_j = \emptyset | X_i = 1] + \mathbb{P}[\mathcal{H}_i \cap \mathcal{H}_j \neq \emptyset | X_i = 1]. \quad (4.100)$$

The result follows from Lemmas 4.15 and 4.16.  $\square$

**Lemma 4.15.** *Let  $i$  and  $j$  be the indices of two vertices on a vertex transitive graph. Then*

$$\mathbb{P}[X_j = 1, \mathcal{H}_i \cap \mathcal{H}_j = \emptyset | X_i = 1] \leq m_{t^*} \quad (4.101)$$

*Proof.* To begin

$$\mathbb{P}[X_j = 1, \mathcal{H}_i \cap \mathcal{H}_j = \emptyset | X_i = 1] = \mathbb{P}[X_i = 1]^{-1} \mathbb{P}[X_i = 1, X_j = 1, \mathcal{H}_i \cap \mathcal{H}_j = \emptyset] \quad (4.102)$$

$$= \frac{\mathbb{P}[\mathcal{H}_i(0) \cup \mathcal{H}_j(0) \neq \emptyset, \mathcal{H}_i \cap \mathcal{H}_j = \emptyset]}{\mathbb{P}[\mathcal{H}_i(0) \neq \emptyset]}. \quad (4.103)$$

Define  $S$  to be the first time that the histories intersect, or define  $S = 0$  if the histories do not intersect before time 0. We construct a new history,  $\mathcal{H}'_j$ , in the following way. Along the interval  $[S, t^*)$ ,  $\mathcal{H}'_j$  is constructed from the update sequence in the exact same way as  $\mathcal{H}_j$ . However, along the interval  $[0, S)$ , we replace the update sequence with an another i.i.d. copy of the update sequence. We have constructed  $\mathcal{H}'_j$  to be independent of  $\mathcal{H}_i$ . This is obvious for  $t < S$ . For  $t \geq S$  we note that no updates are a part of both histories. From the memoryless property of the exponential waiting times between updates, we get that each update is independent of the rest, and so the histories are independent too. Note that in the event  $\{\mathcal{H}_i \cap \mathcal{H}_j = \emptyset\}$ , we have  $\mathcal{H}_j = \mathcal{H}'_j$  and so

$$\mathbb{P}[\mathcal{H}_i(0) \cup \mathcal{H}_j(0) \neq \emptyset, \mathcal{H}_i \cap \mathcal{H}_j = \emptyset] = \mathbb{P}[\mathcal{H}_i(0) \cup \mathcal{H}'_j(0) \neq \emptyset, \mathcal{H}_i \cap \mathcal{H}'_j = \emptyset] \quad (4.104)$$

$$\leq \mathbb{P}[\mathcal{H}_i(0) \cup \mathcal{H}'_j(0) \neq \emptyset] \quad (4.105)$$

$$= \mathbb{P}[\mathcal{H}_i(0) \neq \emptyset] \mathbb{P}[\mathcal{H}'_j(0) \neq \emptyset] \quad (4.106)$$

since  $\mathcal{H}_i$  and  $\mathcal{H}'_j$  are independent. Since  $\mathbb{P}[\mathcal{H}'_j(0) \neq \emptyset] = \mathbb{P}[\mathcal{H}_j(0) \neq \emptyset] = m_{t^*}$  we get the desired result.  $\square$

This proof is closely based on that in [21].

**Lemma 4.16.** *Let  $i$  and  $j$  be the indices of two vertices separated by distance  $k$ . Then there exists a  $C$  such that for sufficiently small  $\beta$ ,*

$$\mathbb{P}[\mathcal{H}_i \cap \mathcal{H}_j \neq \emptyset | X_i = 1] \leq C e^{-k}. \quad (4.107)$$

*Proof.* Firstly,

$$\mathbb{P}[\mathcal{H}_i \cap \mathcal{H}_j \neq \emptyset | X_i = 1] = \frac{\mathbb{P}[\{\mathcal{H}_i \cap \mathcal{H}_j \neq \emptyset\} \cap \{X_i = 1\}]}{\mathbb{P}[X_i = 1]} \quad (4.108)$$

$$\leq m_{t^*}^{-1} \mathbb{P}[\{\mathcal{H}_i \cap \mathcal{H}_j \neq \emptyset\} \cap \{X_i = 1\}] \quad (4.109)$$

$$\leq m_{t^*}^{-1} \mathbb{P}[\{\mathcal{H}_i \cap \mathcal{H}_j \neq \emptyset\} \cap (\{X_i = 1\} \cup \{X_j = 1\})] \quad (4.110)$$

$$= m_{t^*}^{-1} \mathbb{P}[\{\mathcal{H}_i \cap \mathcal{H}_j \neq \emptyset\} \cap \{\mathcal{H}_i(0) \cup \mathcal{H}_j(0) \neq \emptyset\}] \quad (4.111)$$

since

$$\{X_i = 1\} = \{\mathcal{H}_i(0) \neq \emptyset\}. \quad (4.112)$$

Proceeding backwards from  $t^*$ , define  $S$  to be the random time at which  $\mathcal{H}_i(t) \cup \mathcal{H}_j(t)$  first reduced to less than two vertices, or define  $S = 0$  if the combined histories contain at least two vertices all the way to time 0. Note that

$$\{\mathcal{H}_i(0) \cup \mathcal{H}_j(0) \neq \emptyset\} \subseteq \{\mathcal{F}(v, 0, S) \neq \emptyset\} \quad (4.113)$$

where  $v$  is either the single vertex  $v = \mathcal{H}_i(S) \cup \mathcal{H}_j(S)$  in the case that the histories coalesce to a single point at  $S$ , or any arbitrary single vertex otherwise. We also note that

$$\{\mathcal{H}_i \cap \mathcal{H}_j \neq \emptyset\} \subseteq \{\chi((\mathcal{H}_i \cup \mathcal{H}_j) \cap V \times [S, t^*]) \geq k-1\} \quad (4.114)$$

since there must be at least  $k-1$  branching edges for the histories to meet. The event on the right hand side of (4.114) depends only on the update sequence in  $[S, t^*]$ . The event on the right hand side of (4.113) depends only on the update sequence in  $[0, S]$ . Therefore, given  $S$ , these events are independent and

$$\begin{aligned} \mathbb{P}[\{\mathcal{H}_i \cap \mathcal{H}_j \neq \emptyset\} \cap \{\mathcal{H}_i(0) \cup \mathcal{H}_j(0) \neq \emptyset\} | S = t_s] \\ \leq \mathbb{P}[\mathcal{F}(v, 0, S) \neq \emptyset | S = t_s] \mathbb{P}[\chi((\mathcal{H}_i \cup \mathcal{H}_j) \cap V \times [S, t^*]) \geq k-1 | S = t_s] \end{aligned} \quad (4.115)$$

$$= m_{t_s} \mathbb{P}[\chi((\mathcal{H}_i \cup \mathcal{H}_j) \cap V \times [S, t^*]) \geq k-1 | S = t_s] \quad (4.116)$$

and so

$$\mathbb{P}[\{\mathcal{H}_i \cap \mathcal{H}_j \neq \emptyset\} \cap \{\mathcal{H}_i(0) \cup \mathcal{H}_j(0) \neq \emptyset\}] \leq \mathbb{E}[I_{\chi((\mathcal{H}_i \cup \mathcal{H}_j) \cap V \times [S, t^*]) \geq k-1} m_S] \quad (4.117)$$

$$\leq \mathbb{E}[I_{\chi(\mathcal{H}_i \cup \mathcal{H}_j) \geq k-1} m_S]. \quad (4.118)$$

From Corollary 4.3,

$$m_S \leq e^{t^* - S} m_{t^*} \quad (4.119)$$

and since  $|\mathcal{H}_i(t) \cup \mathcal{H}_j(t)| \geq 2$  for  $t \in (S, t^*]$ ,

$$t^* - S \leq \mathcal{L}(\mathcal{H}_i \cup \mathcal{H}_j)/2. \quad (4.120)$$

So

$$\mathbb{P}[\mathcal{H}_i \cap \mathcal{H}_j \neq \emptyset | X_i = 1] \leq m_{t^*}^{-1} m_{t^*} \mathbb{E}[I_{\chi(\mathcal{H}_i \cup \mathcal{H}_j) \geq k-1} e^{\mathcal{L}(\mathcal{H}_i \cup \mathcal{H}_j)/2}] \quad (4.121)$$

$$\leq \mathbb{E}[e^{\chi(\mathcal{H}_i \cup \mathcal{H}_j) - (k-1)} e^{\mathcal{L}(\mathcal{H}_i \cup \mathcal{H}_j)/2}] \quad (4.122)$$

$$= e^{-k+1} \mathbb{E}[e^{\chi(\mathcal{H}_i \cup \mathcal{H}_j) + \mathcal{L}(\mathcal{H}_i \cup \mathcal{H}_j)/2}] \quad (4.123)$$

From Lemma 4.11, for any  $\alpha > \ln(2)$  if

$$\tanh(\beta\Delta) \leq \frac{1 - 2e^{-\alpha}}{2(e^{\Delta(\alpha+1)} - e^{-\alpha})} \quad (4.124)$$

then

$$\mathbb{E}[e^{\chi(\mathcal{H}_i \cup \mathcal{H}_j) + \mathcal{L}(\mathcal{H}_i \cup \mathcal{H}_j)/2}] \leq e^{2\alpha} \quad (4.125)$$

and we get the desired result by choosing  $C = \exp(1 + 2\alpha)$ .

□

## Chapter 5

## Conclusion

## Part II

# Efficient Optimization for Statistical Inference

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