



Computer Graphics (Graphische Datenverarbeitung)

- Splines -

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WS 2021/2022



Corona

- Regular random lookup of the 3G certificates
- Contact tracing: We need to know who is in the class room
 - New ILIAS group for every lecture slot
 - Register via ILIAS or this QR code (only if you are present in this room)





Overview

- Last Time
 - Open-GL
- Today
 - Parametric Curves
 - Lagrange Interpolation
 - Hermite Splines
 - Bézier Splines
 - DeCasteljau Algorithm
 - Parameterization



Curves

Roller Coaster



Roller Coaster – not good



Roller Coaster – not good





Curves

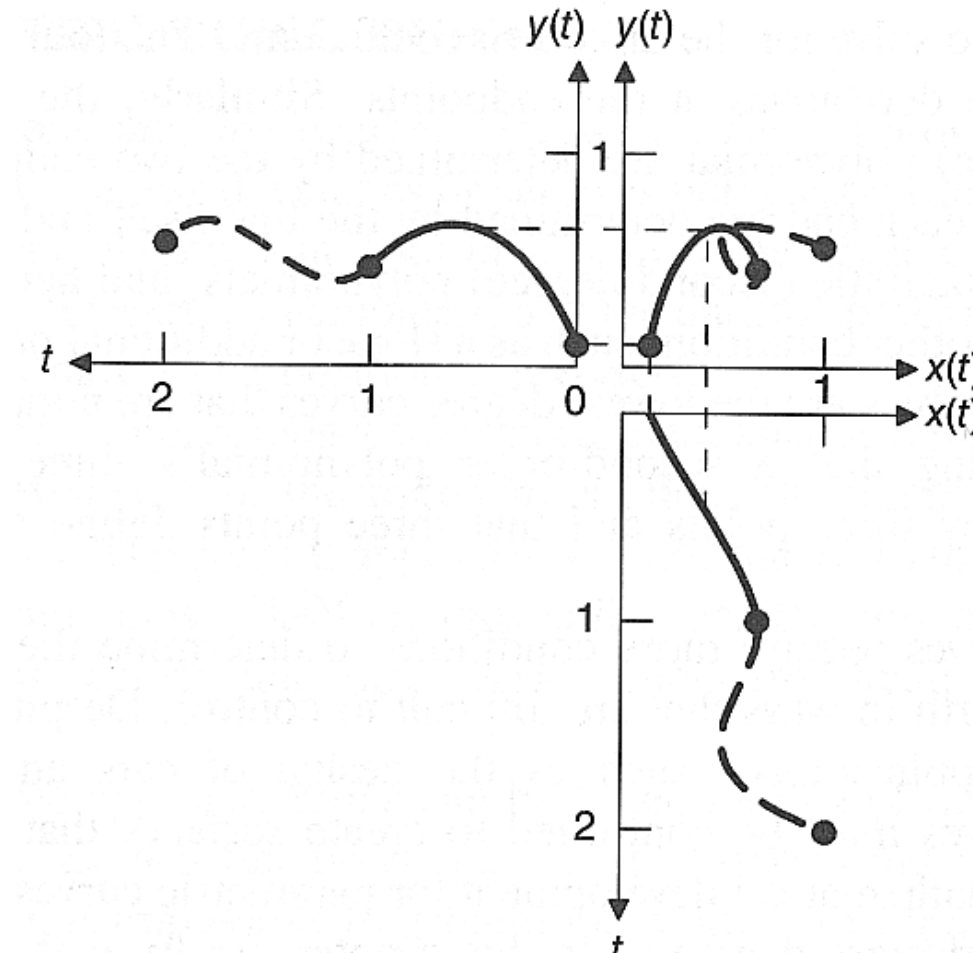
- Curve descriptions
 - Explicit
 - $y(x) = \pm \sqrt{r^2 - x^2}$, restricted domain
 - Implicit:
 - $x^2 + y^2 = r^2$ unknown solution set
 - Parametric:
 - $x(t) = r \cos(t)$, $y(t) = r \sin(t)$, $t \in [0, 2\pi]$
 - Flexibility and ease of use

- Polynomials
 - Avoids complicated functions (z.B. pow, exp, sin, sqrt)
 - Use simple polynomials of low degree



Parametric curves

- Separate function in each coordinate
 - 3D: $f(t) = (x(t), y(t), z(t))$





Monomials

- Monomial basis

- Simple basis: 1, t, t², ... (t usually in [0 .. 1])

$$x(t) = 3t^3 + 1t^2 - 2t + 4$$

- Polynomial representation

$$\underline{P}(t) = \begin{pmatrix} \underline{x}(t) & \underline{y}(t) & \underline{z}(t) \end{pmatrix} = \sum_{i=0}^n t^i \underline{A}_i$$

→ **Degree (= Order – 1)**
→ **Coefficients $\in \mathbb{R}^3$**
→ **Monomials**

- Coefficients can be determined from a sufficient number of constraints (e.g. interpolation of given points)

- Given (n+1) parameter values t_i and points P_i
- Solution of a linear system in the A_i – possible, but inconvenient

- Matrix representation

$$P(t) = \begin{pmatrix} x(t) & y(t) & z(t) \end{pmatrix} = T(t)A = \begin{bmatrix} t^n & t^{n-1} & \dots & 1 \end{bmatrix} \begin{bmatrix} A_{x,n} & A_{y,n} & A_{z,n} \\ A_{x,n-1} & A_{y,n-1} & A_{z,n-1} \\ \vdots & \vdots & \vdots \\ A_{x,0} & A_{y,0} & A_{z,0} \end{bmatrix}$$

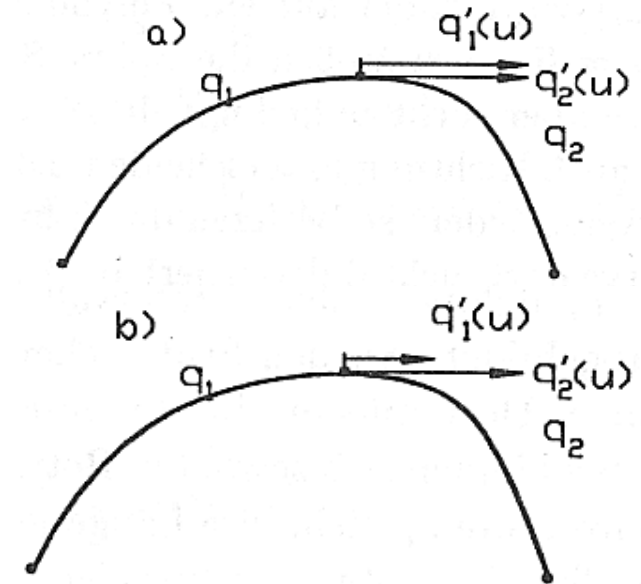


Derivatives

- Derivative = tangent vector
 - Polynomial of degree (n-1)

$$P'(t) = (x'(t) \ y'(t) \ z'(t)) = T'(t)A = [nt^{n-1} \ (n-1)t^{n-2} \ \dots \ 1 \ 0] \begin{bmatrix} A_{x,n} & A_{y,n} & A_{z,n} \\ A_{x,n-1} & A_{y,n-1} & A_{z,n-1} \\ \vdots & \vdots & \vdots \\ A_{x,0} & A_{y,0} & A_{z,0} \end{bmatrix}$$

- Continuity and smoothness between parametric curves
 - $C^0 = G^0$ = same point
 - Parametric continuity C^1
 - Tangent vectors are identical
 - Geometric continuity G^1
 - Same direction of tangent vectors
 - Similar for higher derivatives





More on Continuity

- at one point:
- Geometric Continuity:
 - G^0 : curves are joined
 - G^1 : first derivatives are proportional at joint point, same direction but not necessarily same length
 - G^2 : first and second derivatives are proportional
- Parametric Continuity:
 - C^0 : curves are joined
 - C^1 : first derivative equal
 - C^2 : first and second derivatives are equal. If t is the time, this implies the acceleration is continuous.
 - C^n : all derivatives up to and including the n th are equal.



Lagrange Interpolation

- Interpolating basis functions
 - Lagrange polynomials for a set of parameters $T=\{t_0, \dots, t_n\}$

$$L_i^n(t) = \prod_{\substack{j=0 \\ i \neq j}}^n \frac{t - t_j}{t_i - t_j}, \quad L_i^n(t_j) = \delta_{ij} = \begin{cases} 1 & i = j \\ 0 & i \neq j \end{cases}$$

- Properties
 - Good for interpolation at given parameter values
 - At each t_i : One basis function = 1, all others = 0
 - Polynomial of degree n (n factors linear in t)
- Lagrange Curves
 - Use Lagrange Polynomials with point coefficients

$$\underline{P}(t) = \sum_{i=0}^n L_i^n(t) \underline{P}_i$$



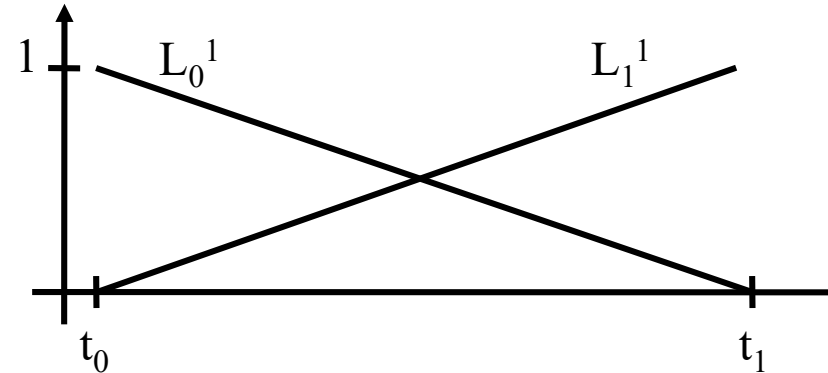
Lagrange Interpolation

- Simple Linear Interpolation

- $T = \{t_0, t_1\}$

$$L_0^1(t) = \frac{t - t_1}{t_0 - t_1}$$

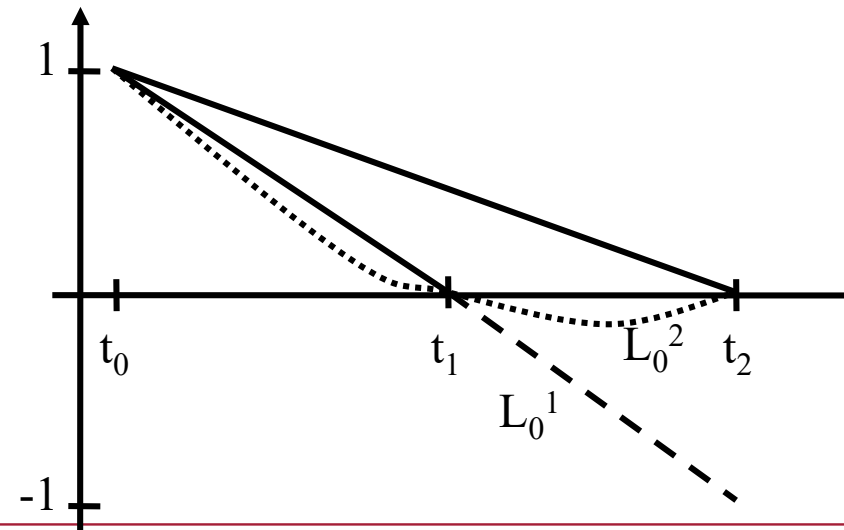
$$L_1^1(t) = \frac{t - t_0}{t_1 - t_0}$$



- Simple Quadratic Interpolation

- $T = \{t_0, t_1, t_2\}$

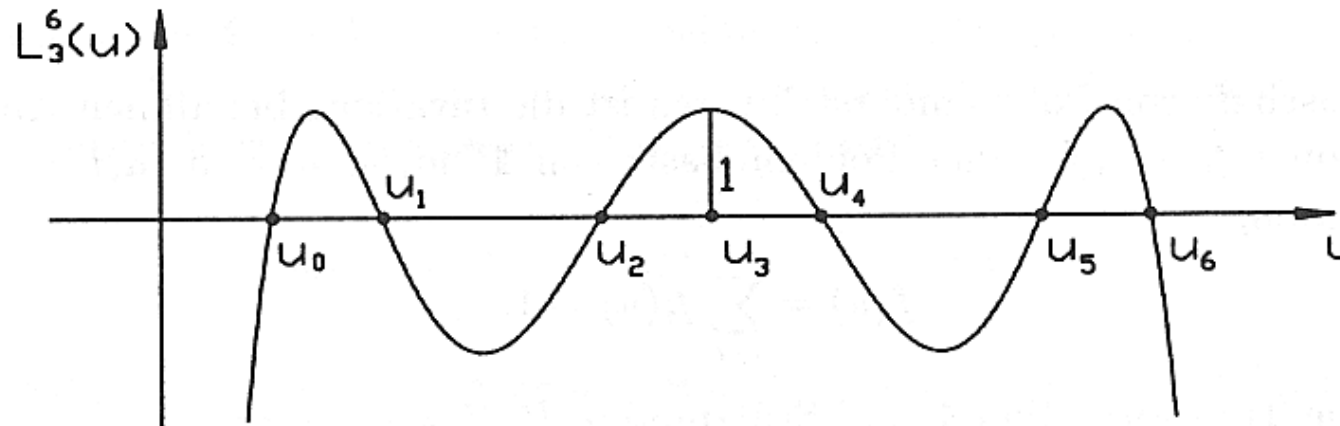
$$L_0^2(t) = \frac{t - t_1}{t_0 - t_1} \frac{t - t_2}{t_0 - t_2}$$





Problems

- Problems with a single polynomial
 - Degree depends on the number of interpolation constraints
 - Strong overshooting for high degree ($n > 7$)
 - Problems with smooth joints
 - Numerically unstable
 - No local changes



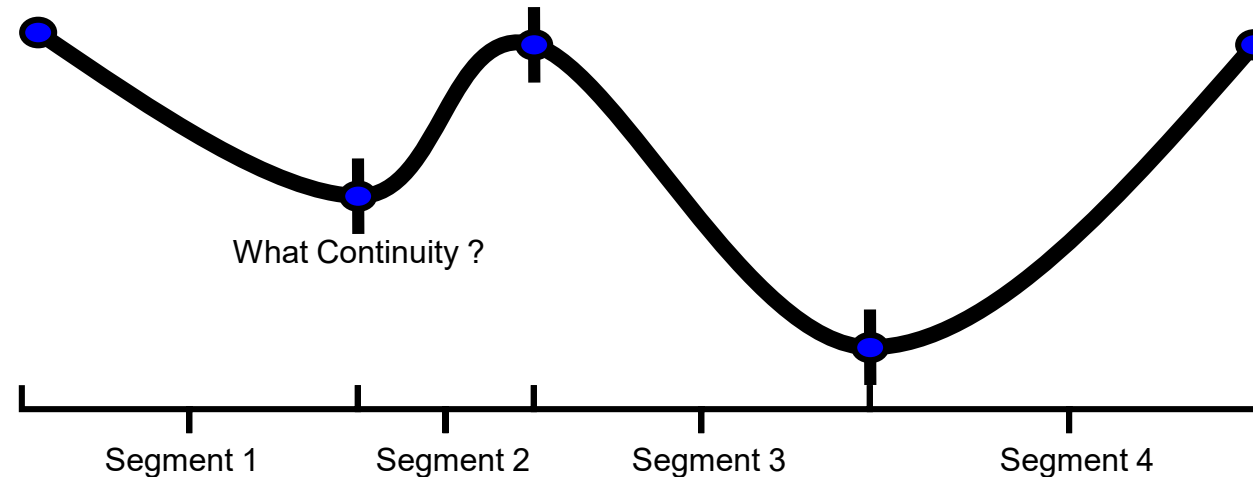


Splines



Splines

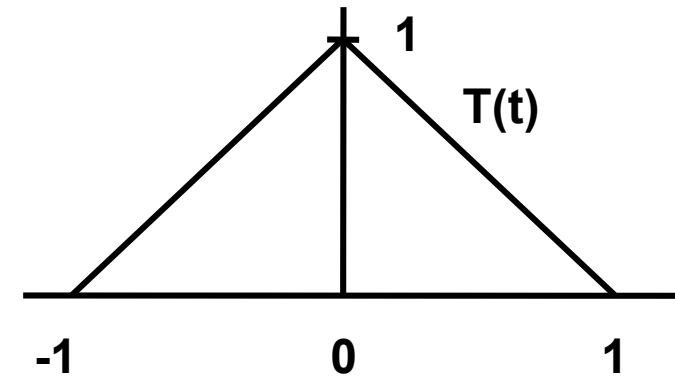
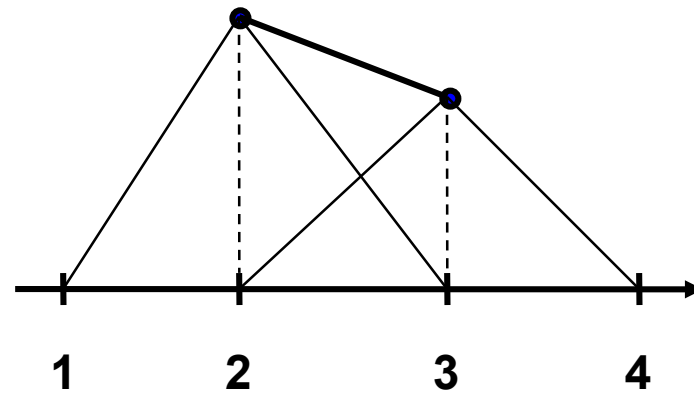
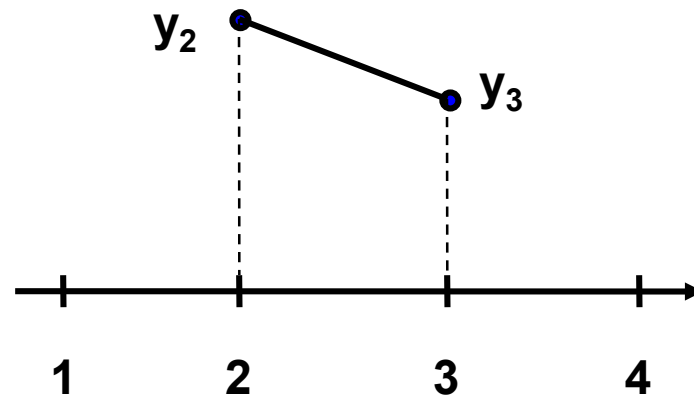
- Functions for interpolation & approximation
 - Standard curve and surface primitives in geometric modeling
 - Key frame and in-betweens in animations
 - Filtering and reconstruction of images
- Historically
 - Name for a tool in ship building
 - Flexible metal strip that tries to stay straight
 - Within computer graphics:
 - Piecewise polynomial function





Linear Interpolation

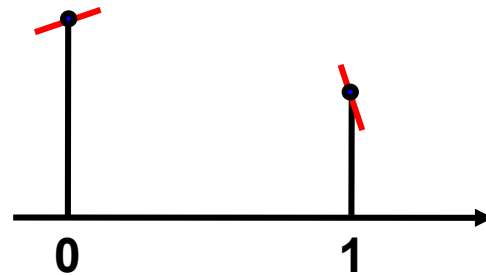
- Hat functions and linear splines
- Piecewise linear function



$$P(t) = T_2(t)y_2 + T_3(t)y_3$$

Hermite Interpolation

- Hermite Basis (cubic)
 - Interpolation of position P and tangent P' information for $t = \{0, 1\}$



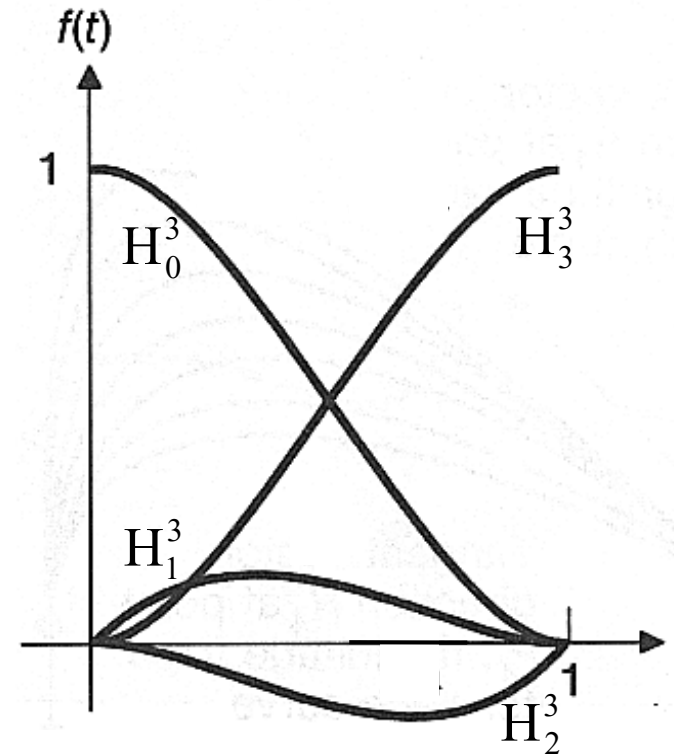
- Basis functions

$$H_0^3(t) = (1-t)^2(1+2t)$$

$$H_1^3(t) = t(1-t)^2$$

$$H_2^3(t) = -t^2(1-t)$$

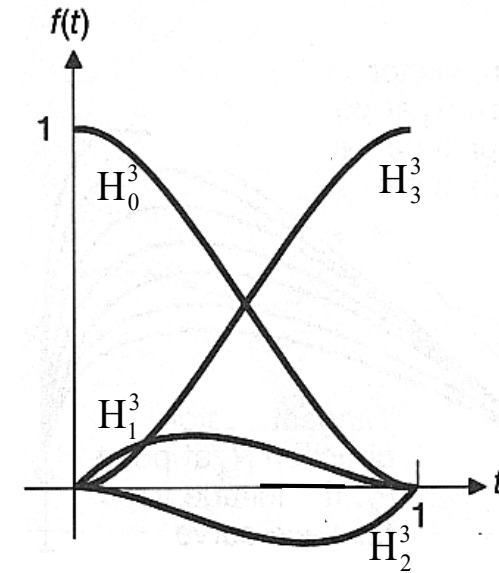
$$H_3^3(t) = (3-2t)t^2$$





Hermite Interpolation

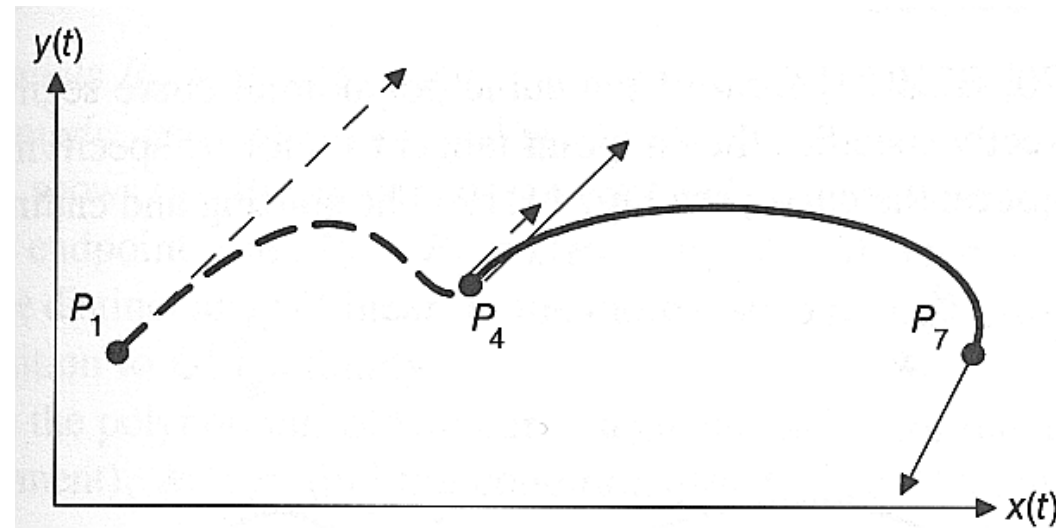
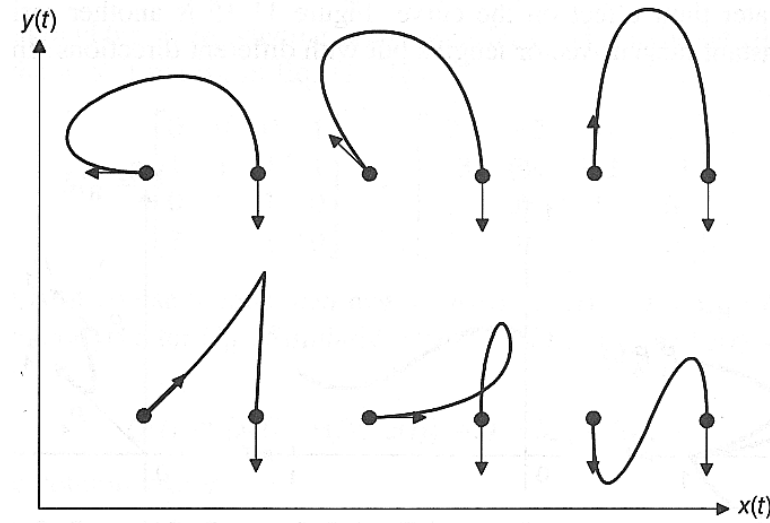
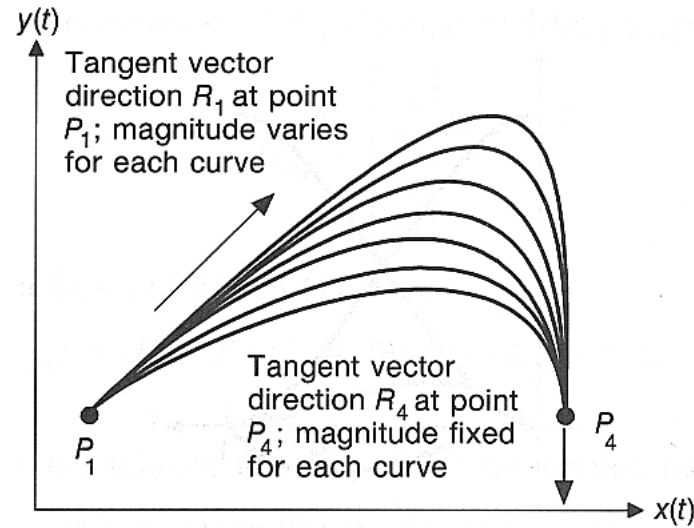
- Properties of Hermite Basis Functions
 - H_0 (H_3) interpolates smoothly from 1 to 0 (1 to 0)
 - H_0 and H_3 have zero derivative at $t=0$ and $t=1$
 - No contribution to derivative (H_1, H_2)
 - H_1 and H_2 are zero at $t=0$ and $t=1$
 - No contribution to position (H_0, H_3)
 - H_1 (H_2) has slope 1 at $t=0$ ($t=1$)
 - Unit factor for specified derivative vector
- Hermite polynomials
 - P_0, P_1 are positions $\in \mathbb{R}^3$
 - P'_0, P'_1 are derivatives (tangent vectors) $\in \mathbb{R}^3$



$$\underline{P}(t) = P_0 H_0^3(t) + P'_0 H_1^3(t) + P'_1 H_2^3(t) + P_1 H_3^3(t)$$



Examples: Hermite Interpolation





Matrix Representation

- Matrix representation

$$\begin{aligned}
 P(t) &= [t^3 \quad t^2 \quad t \quad 1] \begin{bmatrix} A_{x,n} & A_{y,n} & A_{z,n} \\ A_{x,n-1} & A_{y,n-1} & A_{z,n-1} \\ \vdots & \vdots & \vdots \\ A_{x,0} & A_{y,0} & A_{z,0} \end{bmatrix} \\
 &= [t^3 \quad t^2 \quad t \quad 1] \underbrace{\begin{bmatrix} M_{11} & M_{12} & M_{13} & M_{14} \\ M_{21} & \ddots & & \end{bmatrix}}_{\text{basis matrix } M(4 \times 4)} \underbrace{\begin{bmatrix} G_{x,3} & G_{y,3} & G_{z,3} \\ G_{x,2} & G_{y,2} & G_{z,2} \\ G_{x,1} & G_{y,1} & G_{z,1} \\ G_{x,0} & G_{y,0} & G_{z,0} \end{bmatrix}}_{\text{geometry matrix } G(4 \times 3)} \\
 &= [t^3 \quad t^2 \quad t \quad 1] \underbrace{\begin{bmatrix} M_{11} & M_{12} & M_{13} & M_{14} \\ M_{21} & \ddots & & \end{bmatrix}}_{\text{basis functions}} \begin{bmatrix} P_0^T \\ P_1^T \\ P_2^T \\ P_3^T \end{bmatrix}
 \end{aligned}$$



Matrix Representation

- For cubic Hermite interpolation we obtain by evaluating $[t^3 \ t^2 \ t^1 \ t^0]$ or $[3t^2 \ 2t^1 \ 1t^0 \ 0]$ (derivative):

$$\begin{aligned} P_0^T &= (0 \ 0 \ 0 \ 1) \mathbf{M}_H \mathbf{G}_H \\ P_1^T &= (1 \ 1 \ 1 \ 1) \mathbf{M}_H \mathbf{G}_H \\ P_0'^T &= (0 \ 0 \ 1 \ 0) \mathbf{M}_H \mathbf{G}_H \\ P_1'^T &= (3 \ 2 \ 1 \ 0) \mathbf{M}_H \mathbf{G}_H \end{aligned} \quad \text{or} \quad \begin{pmatrix} P_0^T \\ P_1^T \\ P_0'^T \\ P_1'^T \end{pmatrix} = \mathbf{G}_H = \begin{pmatrix} 0 & 0 & 0 & 1 \\ 1 & 1 & 1 & 1 \\ 0 & 0 & 1 & 0 \\ 3 & 2 & 1 & 0 \end{pmatrix} \mathbf{M}_H \mathbf{G}_H$$

- Solution:
 - Two matrices must multiply to identity

$$\mathbf{M}_H = \begin{pmatrix} 0 & 0 & 0 & 1 \\ 1 & 1 & 1 & 1 \\ 0 & 0 & 1 & 0 \\ 3 & 2 & 1 & 0 \end{pmatrix}^{-1} = \begin{pmatrix} 2 & -2 & 1 & 1 \\ -3 & 3 & -2 & -1 \\ 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 \end{pmatrix}$$



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$$\begin{aligned} H_0^3(t) &= (1-t)^2(1+2t) \\ &= 2t^3 - 3t^2 + 1 \end{aligned}$$

$$H_1^3(t) = t(1-t)^2$$

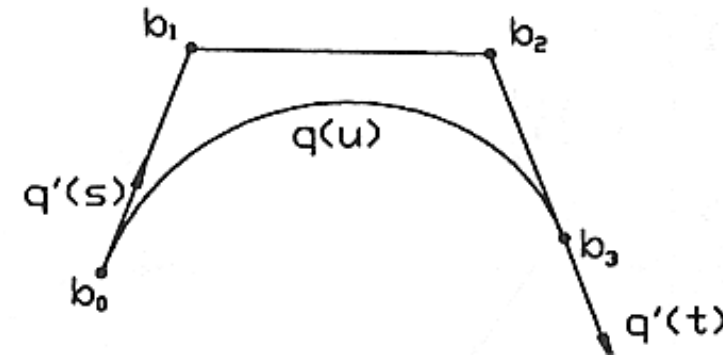
$$H_2^3(t) = -t^2(1-t)$$

$$H_3^3(t) = (3-2t)t^2$$

Bézier

- Bézier Basis [deCasteljau'59, Bézier'62]

- Different curve representation
- Start and end point
- 2 point that are approximated by the curve (cubics)
- $P'_0 = 3(b_1 - b_0)$ and $P'_1 = 3(b_3 - b_2)$
 - Factor 3 due to derivative of t^3



$$G_H = \begin{bmatrix} P_0^T \\ P_1^T \\ P'_0{}^T \\ P'_1{}^T \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ -3 & 3 & 0 & 0 \\ 0 & 0 & -3 & 3 \end{bmatrix} \begin{bmatrix} b_0^T \\ b_1^T \\ b_2^T \\ b_3^T \end{bmatrix} = M_{HB} G_B$$



Basis transformation

- Transformation - Bézier to Hermite

$$- P(t) = T M_H G_H = T M_H (M_{HB} G_B) = T (M_H M_{HB}) G_B = T M_B G_B$$

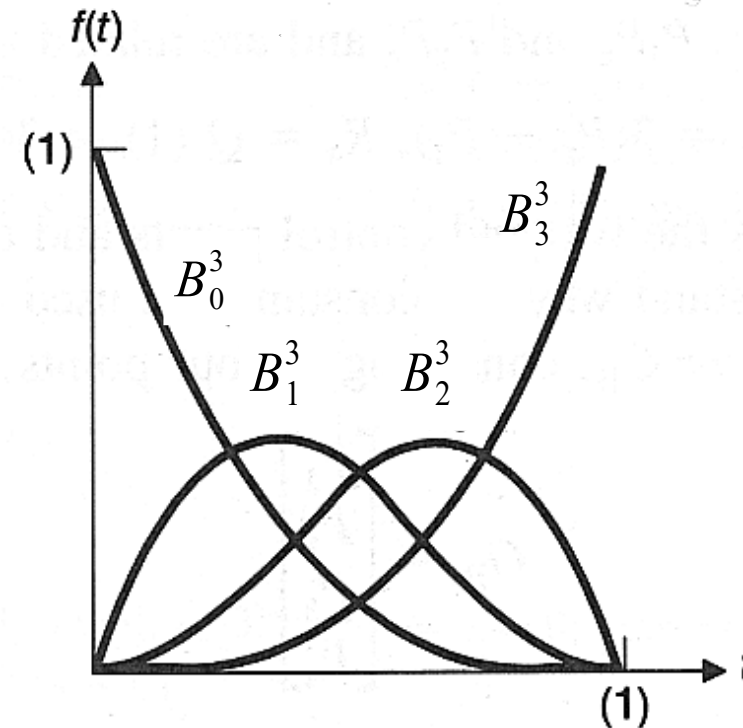
$$M_B = M_H M_{HB} = \begin{bmatrix} -1 & 3 & -3 & 1 \\ 3 & -6 & 3 & 0 \\ -3 & 3 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{bmatrix}$$

- Bézier Curves & Basis Functions

$$P = \sum_{i=0}^3 t^i B_i^3(t) b_i = (1-t)^3 b_0 + 3t(1-t)^2 b_1 + 3t^2(1-t)b_2 + t^3 b_3$$

$$P(t) = \sum_{i=0}^n B_i^n(t) b_i$$

$$\text{with Basisfunctions } B_i^n(t) = \binom{n}{i} t^i (1-t)^{n-i}$$



- Basis functions: **Bernstein polynomials**

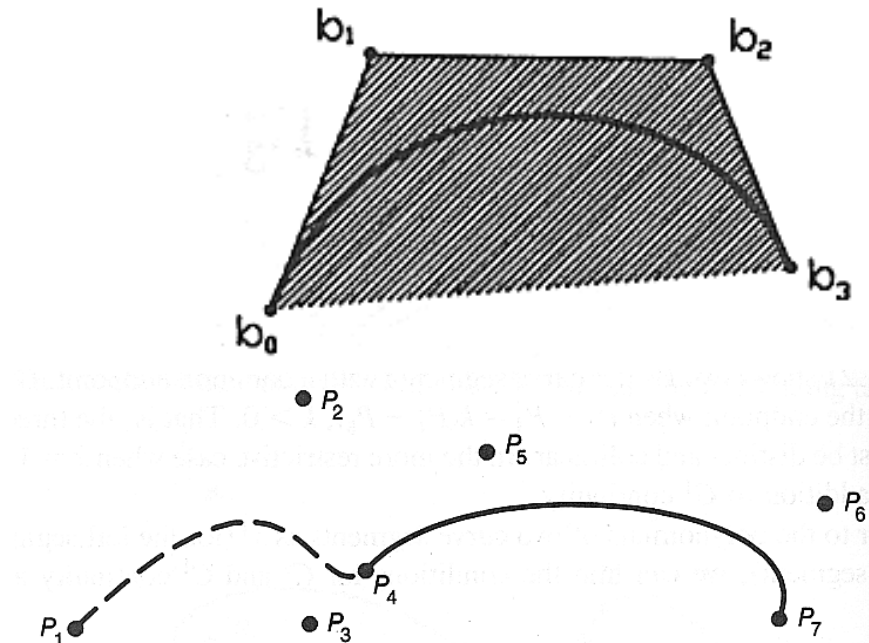
Properties: Bézier

- Advantages:

- End point interpolation
- Tangents explicitly specified
- Smooth joints are simple
 - P_3, P_4, P_5 collinear $\rightarrow G^1$ continuous
- Geometric meaning of control points
- Affine invariance
 - $\forall \sum B_i(t) = 1$
- Convex hull property
 - For $0 < t < 1$: $B_i(t) \geq 0$
- Symmetry: $B_i(t) = B_{n-i}(1-t)$

- Disadvantages

- Smooth joints need to be maintained explicitly
 - Automatic in B-Splines (and NURBS)





DeCasteljau Algorithm



DeCasteljau Algorithm

- Direct evaluation of the basis functions
 - Simple but expensive
- Use recursion
 - Recursive definition of the basis functions

$$B_i^n(t) = tB_{i-1}^{n-1}(t) + (1-t)B_i^{n-1}(t) = \binom{n}{i} t^i (1-t)^{n-i}$$

- Inserting this once yields:

$$P(t) = \sum_{i=0}^n b_i^0 B_i^n(t) = \sum_{i=0}^{n-1} b_i^1(t) B_i^{n-1}(t)$$

- with the new Bézier points given by the recursion

$$b_i^k(t) = t b_{i+1}^{k-1}(t) + (1-t)b_i^{k-1}(t) \quad \text{and} \quad b_i^0(t) = b_i$$

DeCasteljau Algorithm

- DeCasteljau-Algorithm:

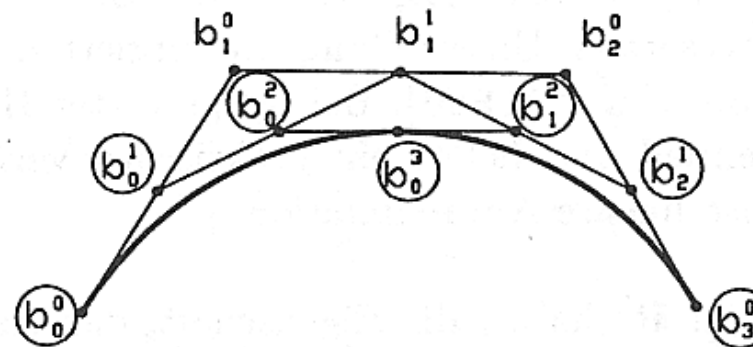
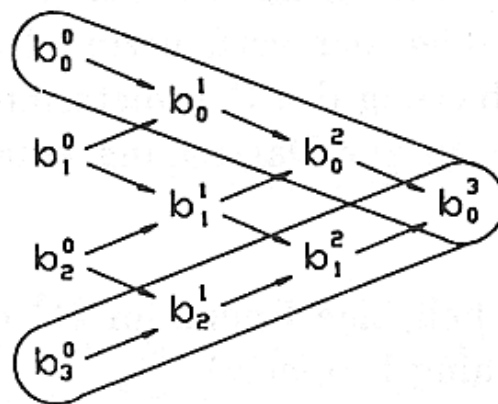
- Recursive degree reduction of the Bezier curve by using the recursion formula for the Bernstein polynomials

$$P(t) = \sum_{i=0}^n b_i^0 B_i^n(t) = \sum_{i=0}^{n-1} b_i^1 B_i^{n-1}(t) = \dots = b_i^n(t) \cdot 1$$

- Example:

- $t = 0.5$

$$b_i^k(t) = t b_{i+1}^{k-1}(t) + (1-t) b_i^{k-1}(t)$$



DeCasteljau Algorithm

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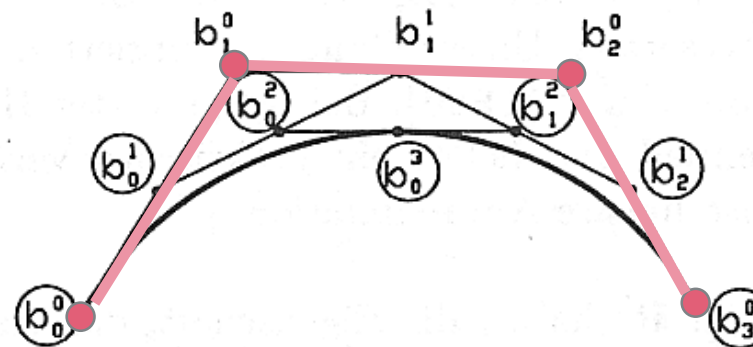
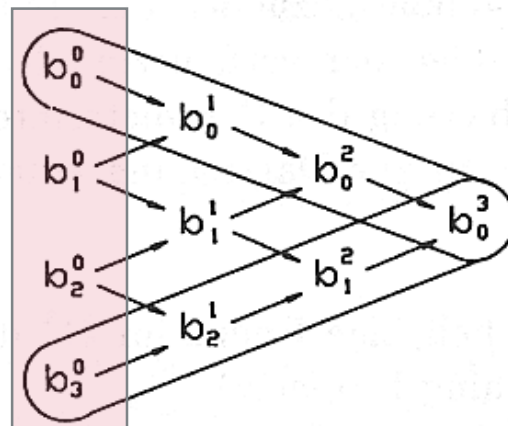
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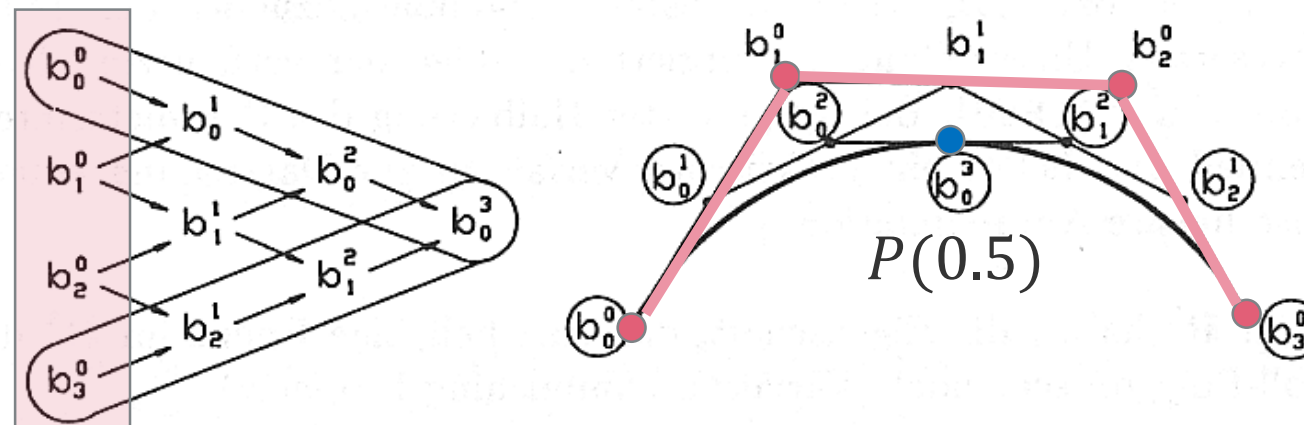
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DeCasteljau Algorithm

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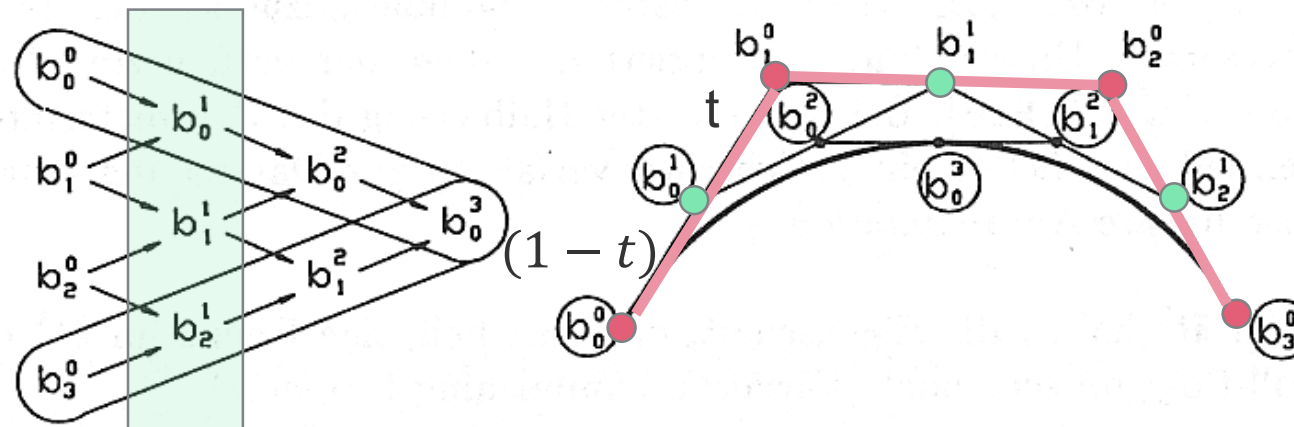
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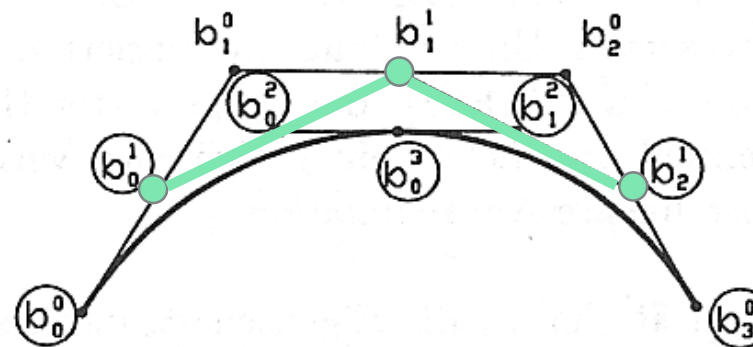
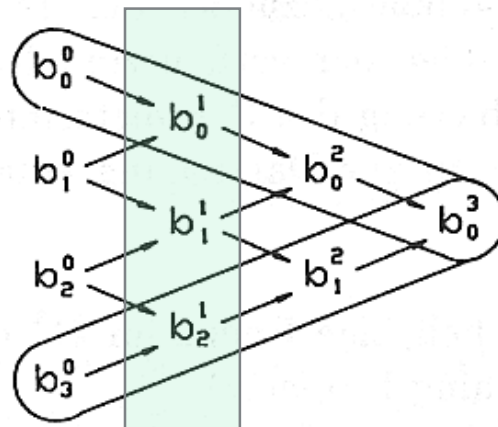
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DeCasteljau Algorithm

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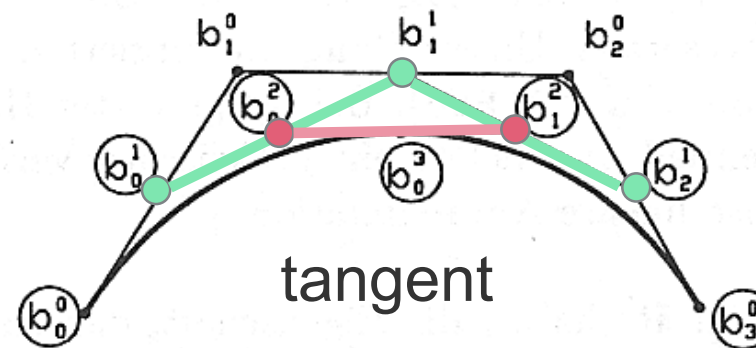
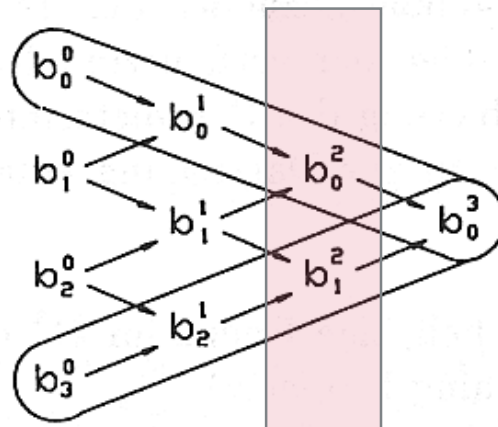
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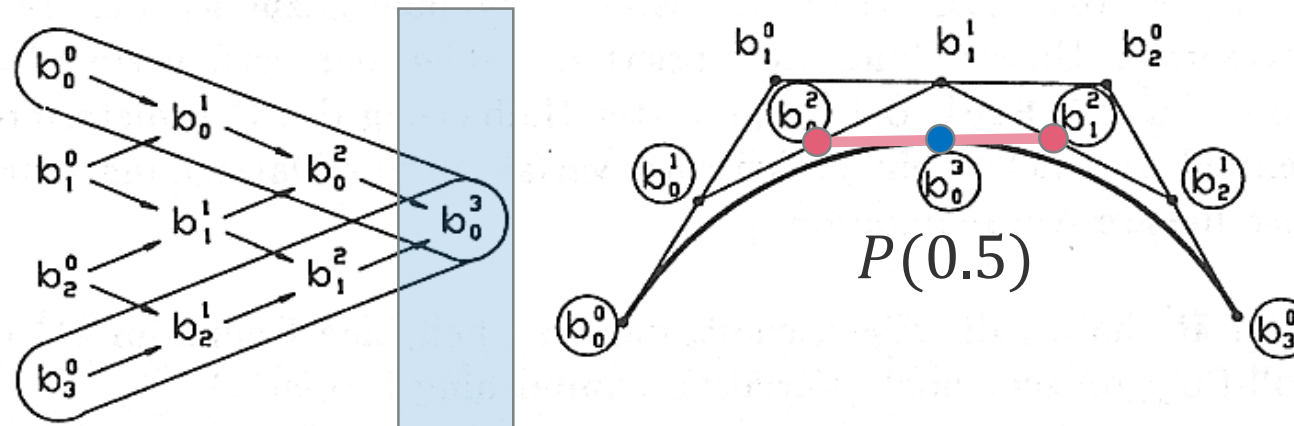
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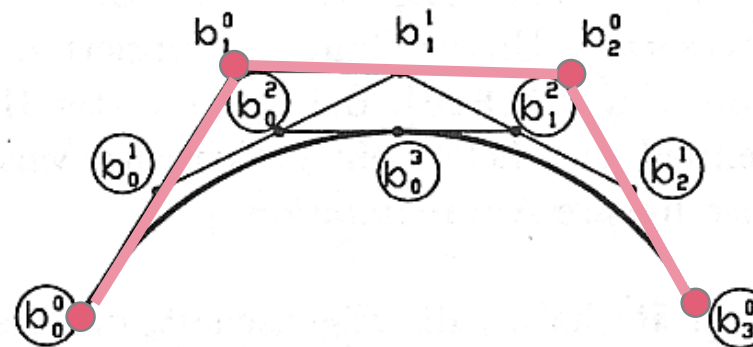
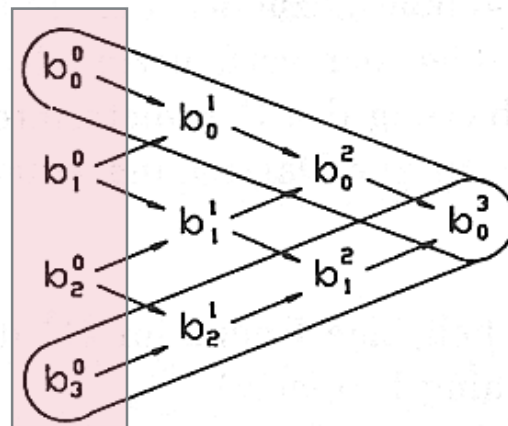
- Recursive degree reduction of the Bezier curve by using the recursion formula for the Bernstein polynomials

$$P(t) = \sum_{i=0}^n b_i^0 B_i^n(t) = \sum_{i=0}^{n-1} b_i^1 B_i^{n-1}(t) = \dots = b_i^n(t) \cdot 1$$

- Example:

- $t = 0.25$

$$b_i^k(t) = t b_{i+1}^{k-1}(t) + (1-t) b_i^{k-1}(t)$$



DeCasteljau Algorithm

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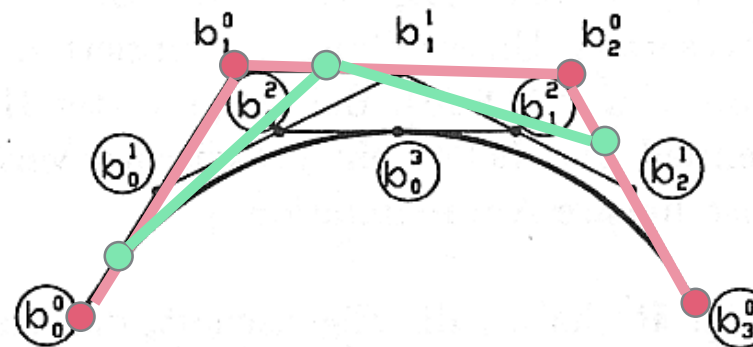
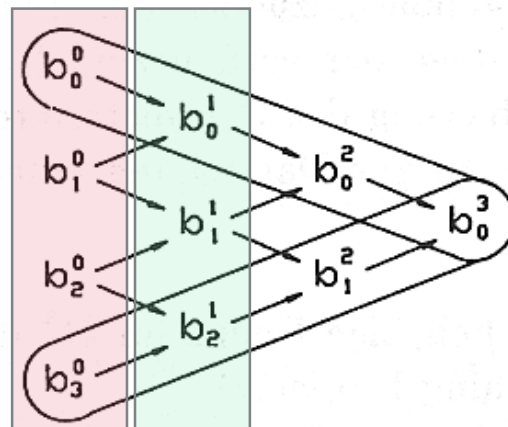
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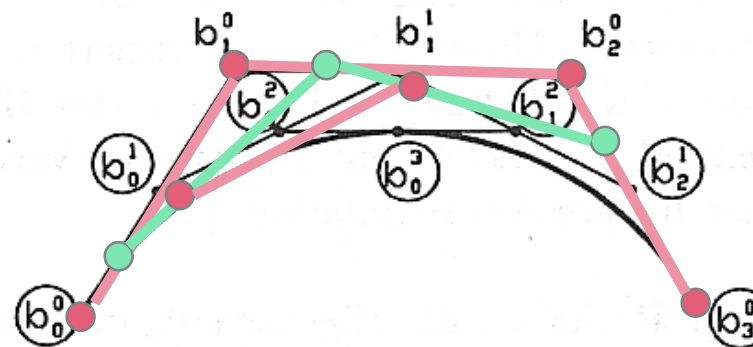
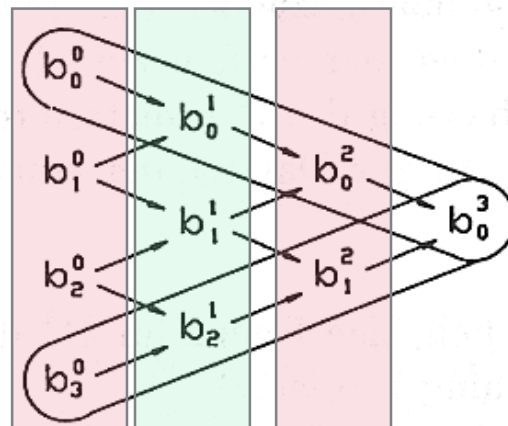
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DeCasteljau Algorithm

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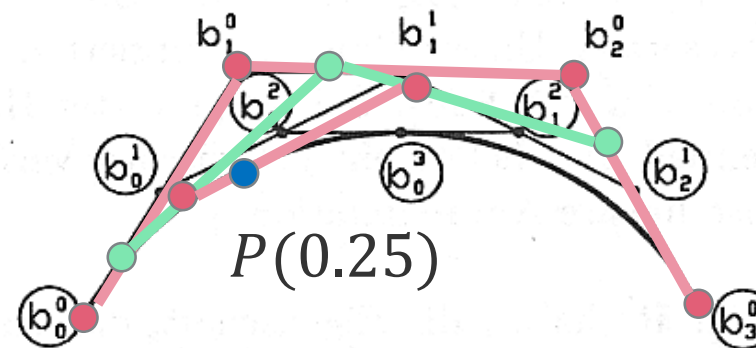
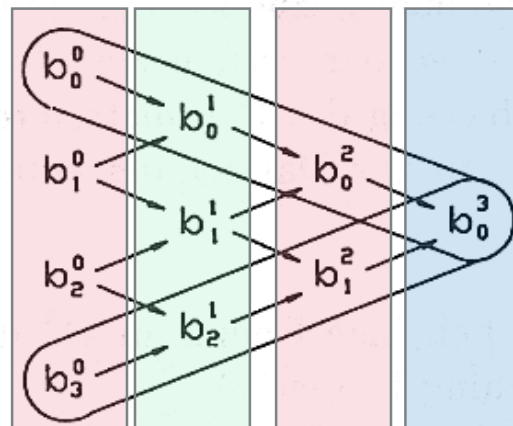
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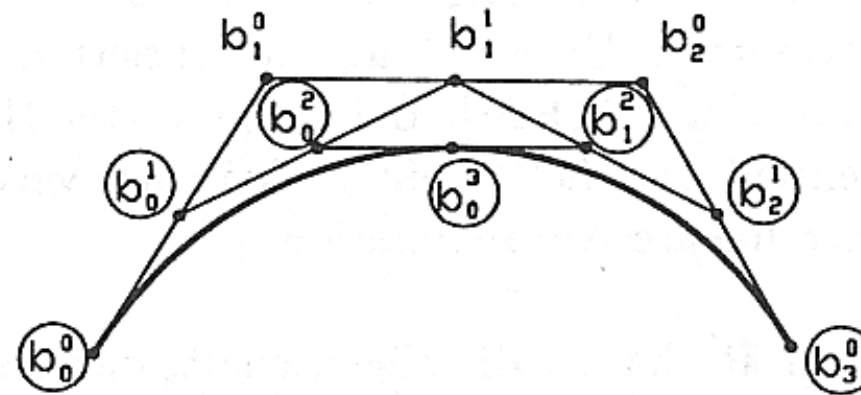
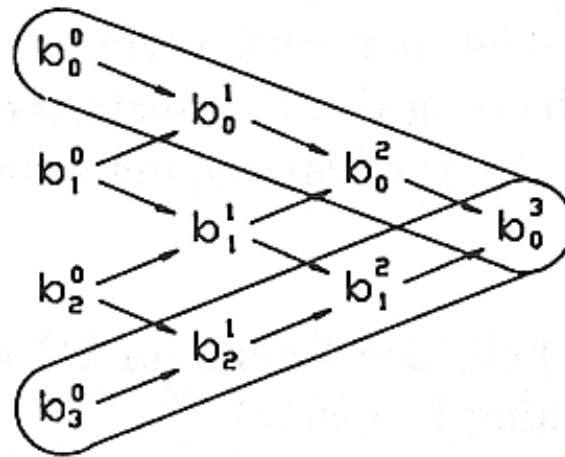
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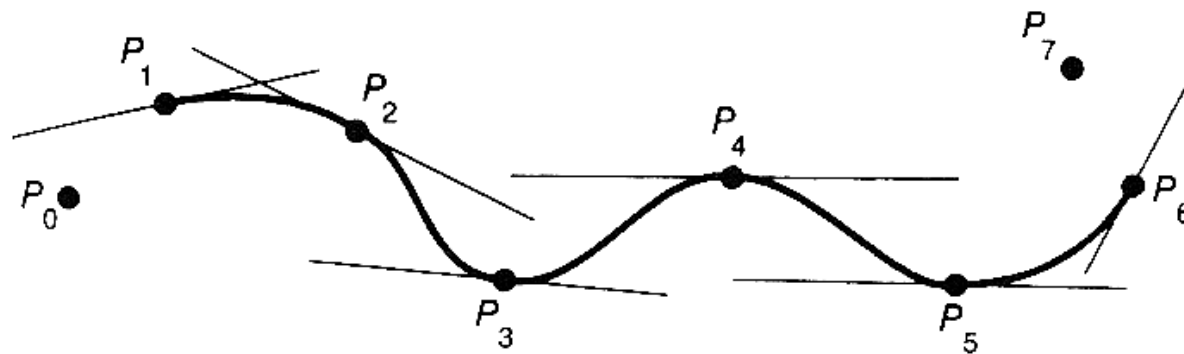
DeCasteljau Algorithm

- Subdivision using the deCasteljau-Algorithm
 - Take boundaries of the deCasteljau triangle as new control points for left/right portion of the curve
- Extrapolation
 - Backwards subdivision
 - Reconstruct triangle from one side



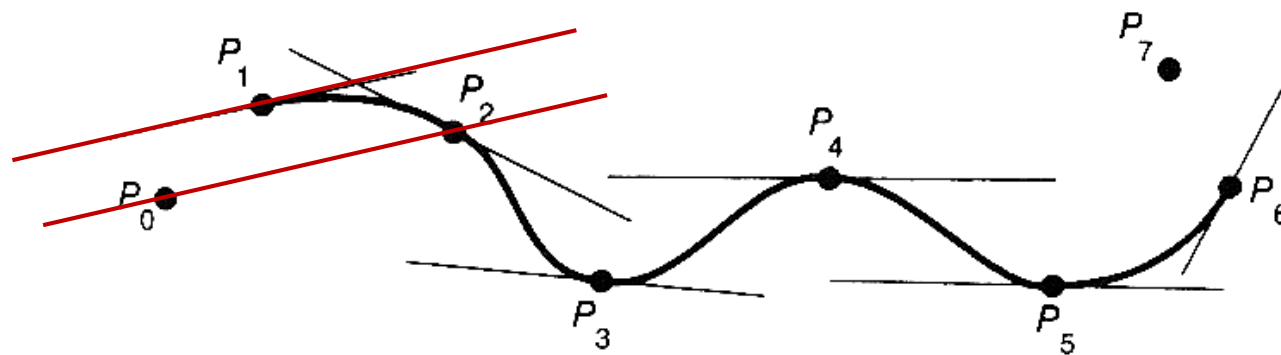
Catmull-Rom-Splines

- Goal
 - Smooth (C^1)-joints between (cubic) spline segments
- Algorithm
 - Tangents given by neighboring points P_{i-1} P_{i+1}
 - Construct (cubic) Hermite segments
- Advantage
 - Arbitrary number of control points
 - Interpolation without overshooting
 - Local control



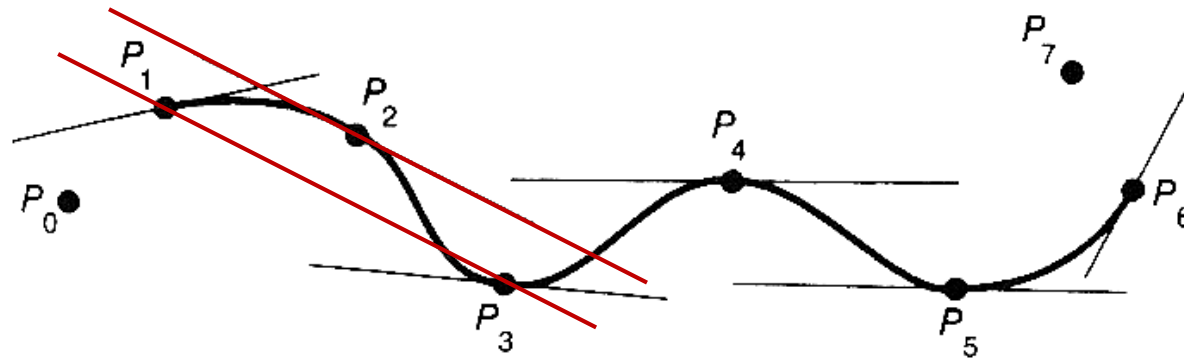
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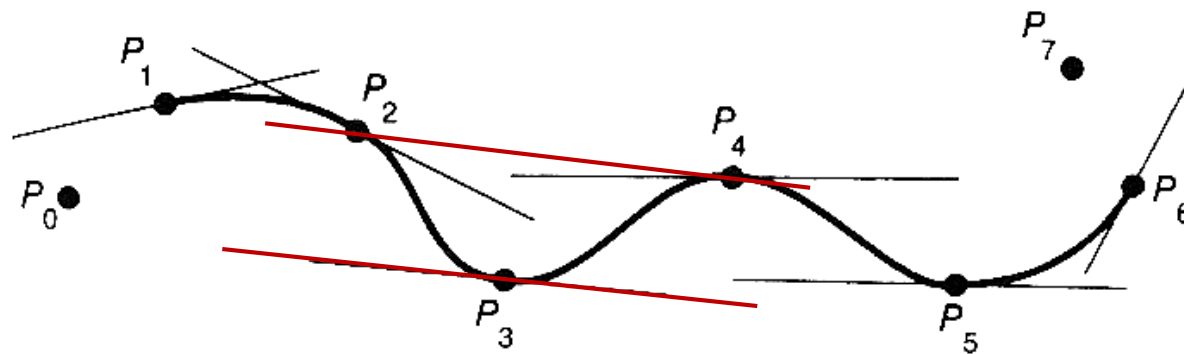
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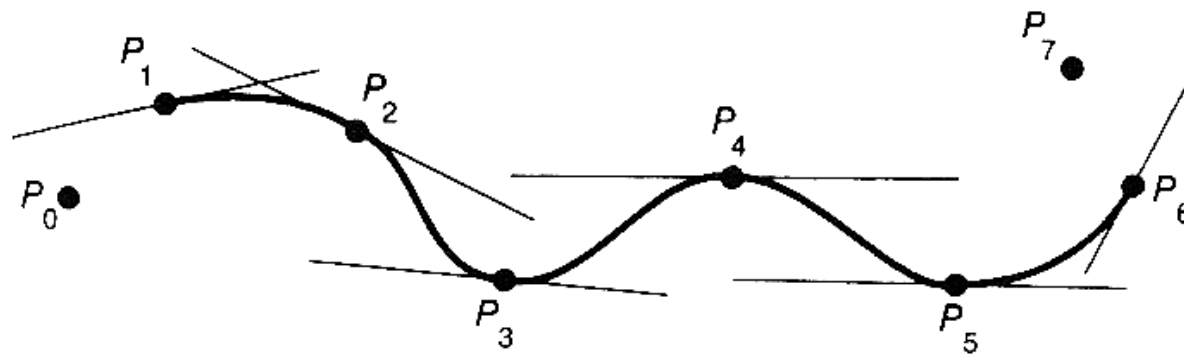
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Matrix Representation

- Catmull-Rom-Spline

- Piecewise polynomial curve
- Four control points per segment
- For n control points we obtain (n-3) polynomial segments

$$\underline{P}^i(t) = T \mathbf{M}_C G_{RC} \bar{R} T \frac{1}{2} \begin{bmatrix} -1 & 3 & -3 & 1 \\ 2 & -5 & 4 & 1 \\ -1 & 0 & 1 & 0 \\ 0 & 2 & 0 & 0 \end{bmatrix} \begin{bmatrix} \underline{P}_i^T \\ \underline{P}_{i+1}^T \\ \underline{P}_{i+2}^T \\ \underline{P}_{i+3}^T \end{bmatrix}$$

- Application

- Smooth interpolation of a given sequence of points
- Key frame animation, camera movement, etc.
- Only G¹-continuity
- Control points should be equidistant in time



B-Splines



Choice of Parameterization

- Problem
 - Often only the control points are given
 - How to obtain a suitable parameterization t_i ?
- Example: Chord-Length Parameterization

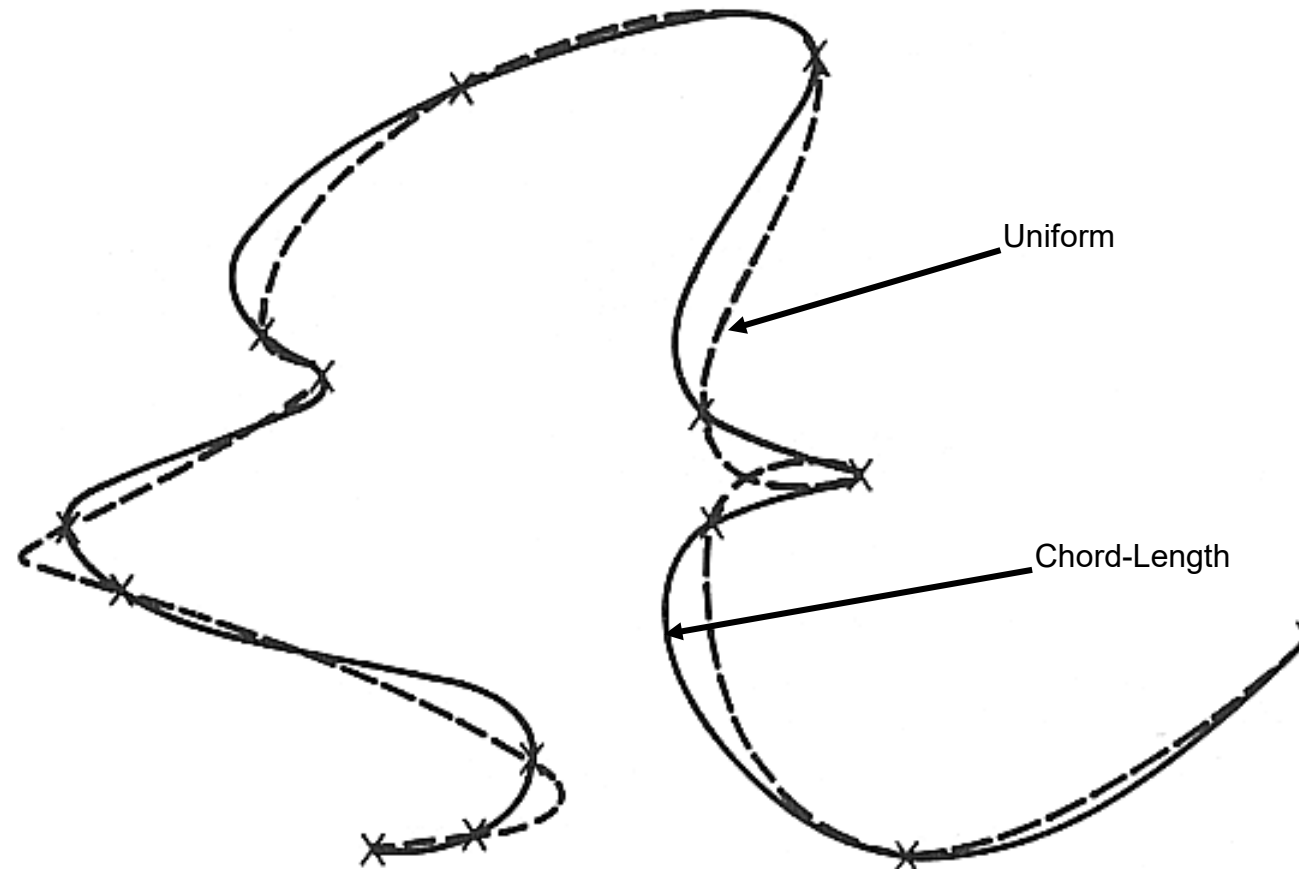
$$t_0 = 0$$

$$t_i = \sum_{j=1}^i d_j \quad d_j = \|P_j - P_{j-1}\|$$

- Arbitrary up to a constant factor
- Warning
 - Distances are not affine invariant !
 - Shape of curves changes under transformations !!

Parameterization

- Chord-Length versus uniform Parameterization
 - Analog: Think $P(t)$ as a moving object with mass that may overshoot





B-Splines

- Goal
 - Spline curve with local control and high continuity
- Given
 - Degree: n
 - Control points: P_0, \dots, P_m (Control polygon, $m \geq n+1$)
 - Knots: t_0, \dots, t_{m+n+1} (Knot vector, weakly monotonic)
 - The knot vector defines the parametric locations where segments join
- B-Spline Curve

$$\underline{P}(t) = \sum_{i=0}^m N_i^n(t) \underline{P}_i$$

- Continuity:
 - C_{n-1} at simple knots
 - C_{n-k} at knot with multiplicity k

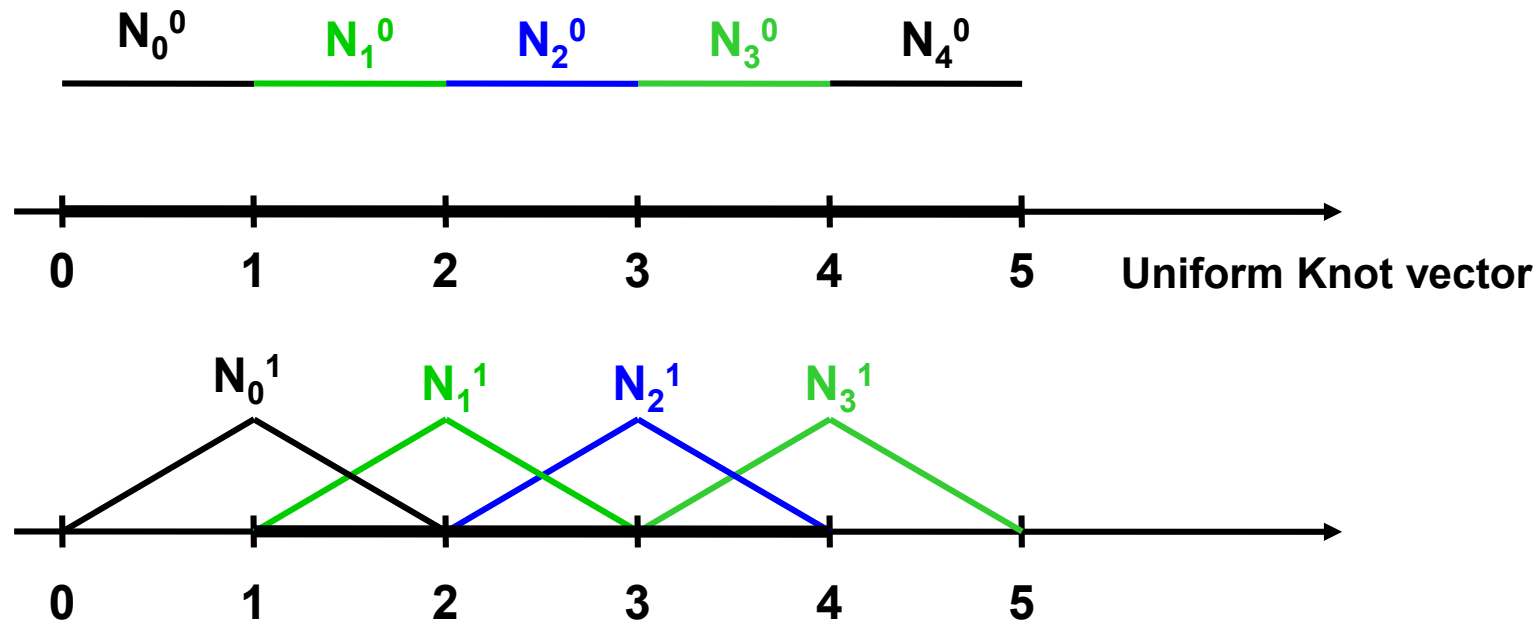


B-Spline Basis Functions

- Recursive Definition

$$N_i^0(t) = \begin{cases} 1 & \text{if } t_i < t < t_{i+1} \\ 0 & \text{otherwise} \end{cases}$$

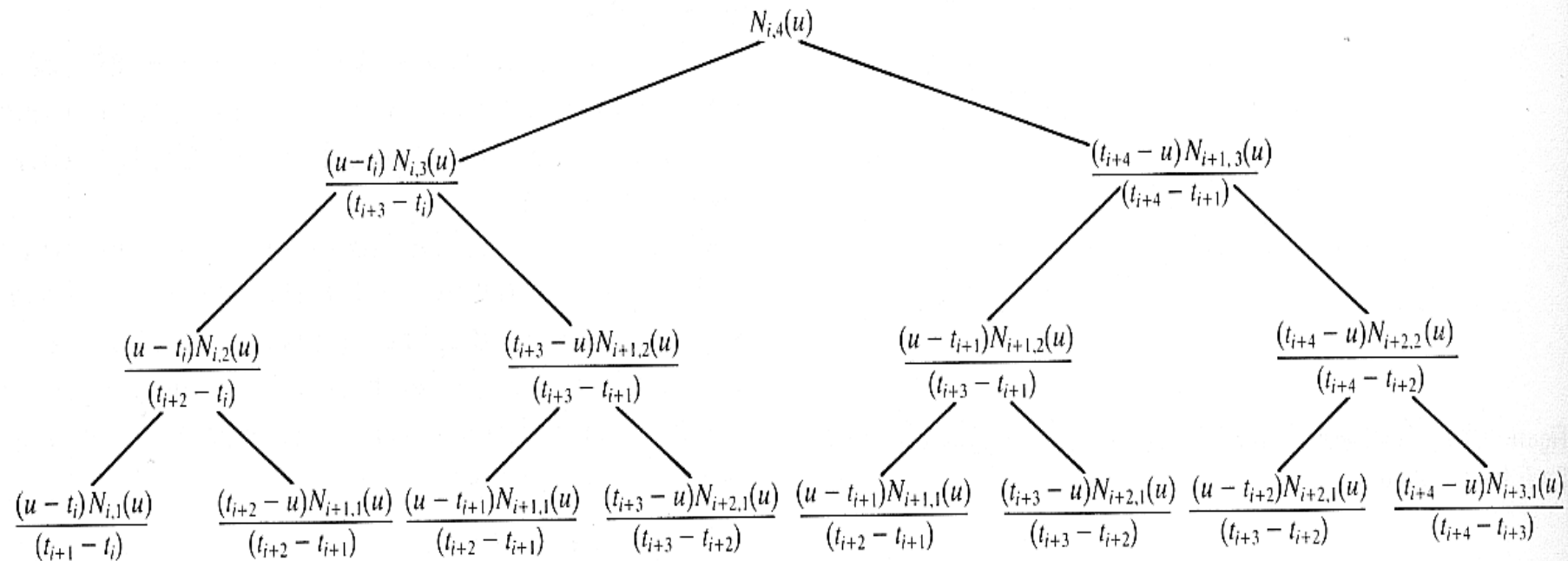
$$N_i^n(t) = \frac{t - t_i}{t_{i+n} - t_i} N_i^{n-1}(t) + \frac{t_{i+1} - t}{t_{i+1} - t_{i+2}} N_{i+1}^{n-1}(t)$$





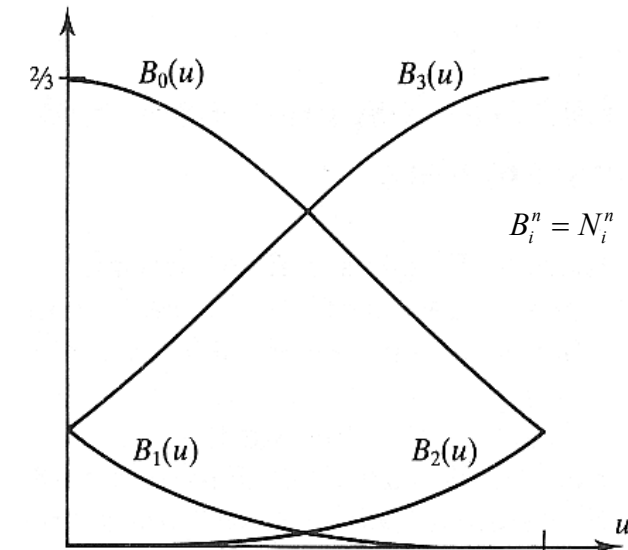
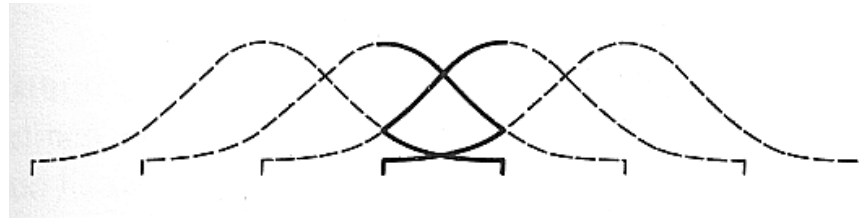
B-Spline Basis Functions

- Recursive Definition
 - Degree increases in every step
 - Support increases by one knot interval

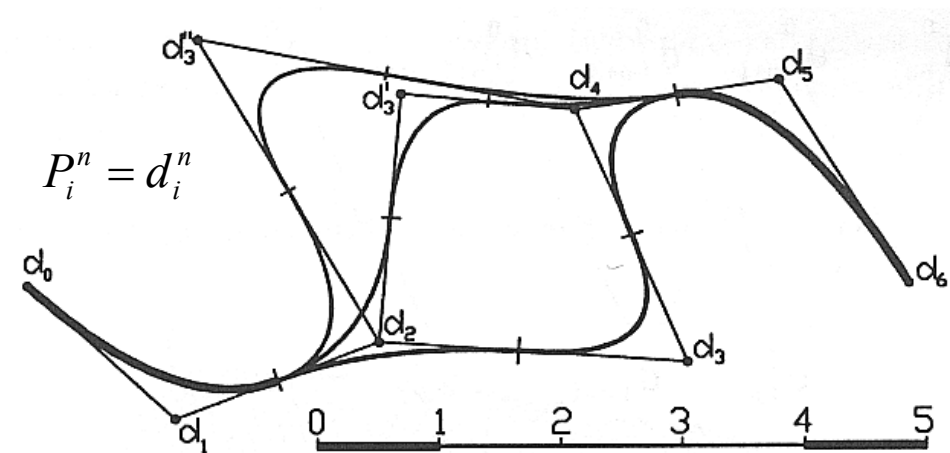


B-Spline Basis Functions

- Uniform Knot Vector
 - All knots at integer locations
 - UBS: Uniform B-Spline
 - Example: cubic B-Splines



- Local Support = Localized Changes
 - Basis functions affect only $(n+1)$ Spline segments
 - Changes are localized



Degree 2



Summary

- Interpolating polynomials hard to control
- Splines: Curves as piece-wise polynomial functions
- Matrices for basis transformations / calculation of derivatives
- DeCasteljau Algorithm for efficient evaluation
- B-Splines can control the parameterization



- Tuesday, 22.02.2022
- 08:00 (s.t.) – 11:00, N10!
- Bring a Pen and a Ruler
- No calculator allowed.
- Resit exam: **Thursday, 31.03.2022!**