

# Chapter 6

## Relations

*Discrete Structures for Computing* on January 4, 2023

Relations

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Toan, Tran Hong Tai



Contents

Properties of Relations

Combining Relations

Representing Relations

Closures of Relations

Types of Relations

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# Course outcomes

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## Course learning outcomes

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L.O.1	Understanding of logic and discrete structures
	L.O.1.1 – Describe definition of propositional and predicate logic
	L.O.1.2 – Define basic discrete structures: set, mapping, graphs
L.O.2	Represent and model practical problems with discrete structures
	L.O.2.1 – Logically describe some problems arising in Computing
	L.O.2.2 – Use proving methods: direct, contrapositive, induction
	L.O.2.3 – Explain problem modeling using discrete structures
L.O.3	Understanding of basic probability and random variables
	L.O.3.1 – Define basic probability theory
	L.O.3.2 – Explain discrete random variables
L.O.4	Compute quantities of discrete structures and probabilities
	L.O.4.1 – Operate (compute/ optimize) on discrete structures
	L.O.4.2 – Compute probabilities of various events, conditional ones, Bayes theorem

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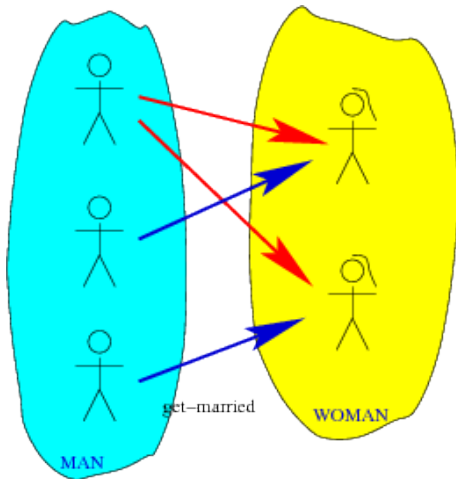
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# Introduction



## Function?

### Relations

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### Definition

Let  $A$  and  $B$  be sets. A **binary relation** (*quan hệ hai ngôi*) from a set  $A$  to a set  $B$  is a set

$$R \subseteq A \times B$$

- Notations:

$$(a, b) \in R \longleftrightarrow aRb$$

- **n-ary relations?**

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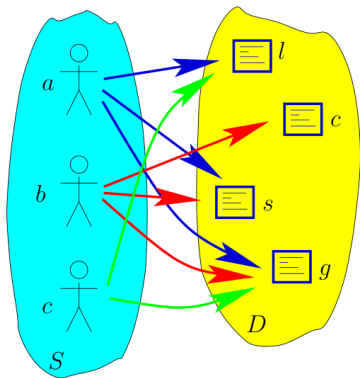
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## Example

### Example

Let  $A = \{a, b, c\}$  be the set of students,  $B = \{l, c, s, d\}$  be the set of the available optional courses. We can have relation  $R$  that consists of pairs  $(a, b)$ , where  $a$  is a student enrolled in course  $b$ .



$$R = \{(a, l), (a, s), (a, g), (b, c), (b, s), (b, g), (c, l), (c, g)\}$$

$R$	$l$	$c$	$s$	$g$
$a$	x		x	x
$b$		x	x	x
$c$	x			x



# Functions as Relations

- Is a function a relation?

- **Yes!**

- $f : A \rightarrow B$

$$R = \{(a, b) \mid b = f(a)\}$$

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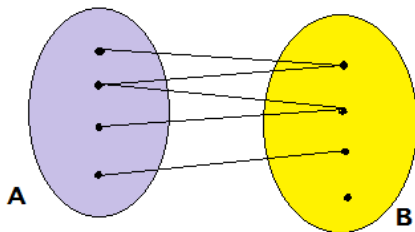
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# Functions as Relations

- Is a relation a function?
- **No**



- Relations are a **generalization** of functions





# Relations on a Set

## Definition

A **relation** on the set  $A$  is a relation from  $A$  to  $A$ .

## Example

Let  $A$  be the set  $\{1, 2, 3, 4\}$ . Which ordered pairs are in the relation  $R = \{(a, b) \mid a \text{ divides } b\}$  ( $a$  là ước số của  $b$ )?

Solution:

$$R = \{(1, 1), (1, 2), (1, 3), (1, 4), (2, 2), (2, 4), (3, 3), (4, 4)\}$$

$R$	1	2	3	4
1	x	x	x	x
2		x		x
3			x	
4				x



# Properties of Relations



<b>Reflexive</b> (phản xạ)	$xRx, \forall x \in A$
<b>Symmetric</b> (đối xứng)	$xRy \rightarrow yRx, \forall x, y \in A$
<b>Antisymmetric</b> (phản đối xứng)	$(xRy \wedge yRx) \rightarrow x = y, \forall x, y \in A$
<b>Transitive</b> (bắc cầu)	$(xRy \wedge yRz) \rightarrow xRz, \forall x, y, z \in A$

## Example

### Example

Consider the following relations on  $\{1, 2, 3, 4\}$ :

$$R_1 = \{(1, 1), (1, 2), (2, 1), (2, 2), (3, 4), (4, 1), (4, 4)\},$$

$$R_2 = \{(1, 1), (1, 2), (2, 1)\},$$

$$R_3 = \{(1, 1), (1, 2), (1, 4), (2, 1), (2, 2), (3, 3), (4, 1), (4, 4)\},$$

$$R_4 = \{(2, 1), (3, 1), (3, 2), (4, 1), (4, 2), (4, 3)\},$$

$$R_5 = \{(3, 4)\}$$

### Solution:

- Reflexive:  $R_3$
- Symmetric:  $R_2, R_3$
- Antisymmetric:  $R_4, R_5$
- Transitive:  $R_4, R_5$



## Example

### Example

What is the properties of the **divides** (ước số) relation on the set of positive integers?

### Solution:

- $\forall a \in \mathbb{Z}^+, a \mid a$ : **reflexive**
- $1 \mid 2$ , but  $2 \nmid 1$ : **not symmetric**
- $\exists a, b \in \mathbb{Z}^+, (a \mid b) \wedge (b \mid a) \rightarrow a = b$ : **antisymmetric**
- $a \mid b \Rightarrow \exists k \in \mathbb{Z}^+, b = ak; b \mid c \Rightarrow \exists l \in \mathbb{Z}^+, c = bl$ . Hence,  $c = a(kl) \Rightarrow a \mid c$ : **transitive**



## Example

### Example

What are the properties of these relations on the set of integers:

$$R_1 = \{(a, b) \mid a \leq b\}$$

$$R_2 = \{(a, b) \mid a > b\}$$

$$R_3 = \{(a, b) \mid a = b \text{ or } a = -b\}$$



## Combining Relations

Because relations from  $A$  to  $B$  are **subsets** of  $A \times B$ , two relations from  $A$  to  $B$  can be combined in any way two sets can be combined.

### Example

Let  $A = \{1, 2, 3\}$  and  $B = \{1, 2, 3, 4\}$ . List the combinations of relations  $R_1 = \{(1, 1), (2, 2), (3, 3)\}$  and  $R_2 = \{(1, 1), (1, 2), (1, 3), (1, 4)\}$ .

**Solution:**  $R_1 \cup R_2$ ,  $R_1 \cap R_2$ ,  $R_1 - R_2$  and  $R_2 - R_1$ .

### Example

Let  $A$  and  $B$  be the set of all students and the set of all courses at school, respectively. Suppose  $R_1 = \{(a, b) \mid a \text{ has taken the course } b\}$  and  $R_2 = \{(a, b) \mid a \text{ requires course } b \text{ to graduate}\}$ . What are the relations  $R_1 \cup R_2$ ,  $R_1 \cap R_2$ ,  $R_1 \oplus R_2$ ,  $R_1 - R_2$ ,  $R_2 - R_1$ ?



# Composition of Relations



## Definition

Let  $R$  be **relations** from  $A$  to  $B$  and  $S$  be from  $B$  to  $C$ . Then the **composite** (*hợp thành*) of  $S$  and  $R$  is

$$S \circ R = \{(a, c) \in A \times C \mid \exists b \in B (aRb \wedge bSc)\}$$

## Example

$$R = \{(0, 0), (0, 3), (1, 2), (0, 1)\}$$

$$S = \{(0, 0), (1, 0), (2, 1), (3, 1)\}$$

$$S \circ R = \{(0, 0), (0, 1), (1, 1)\}$$

# Power of Relations

## Definition

Let  $R$  be a relation on the set  $A$ . The **powers** (*lũy thừa*)  $R^n, n = 1, 2, 3, \dots$  are defined recursively by

$$R^1 = R \quad \text{and} \quad R^{n+1} = R^n \circ R.$$

## Example

Let  $R = \{(1, 1), (2, 1), (3, 2), (4, 3)\}$ . Find the powers  $R^n, n = 2, 3, 4, \dots$

### Solution:

$$R^2 = \{(1, 1), (2, 1), (3, 1), (4, 2)\}$$

$$R^3 = \{(1, 1), (2, 1), (3, 1), (4, 1)\}$$

$$R^4 = \{(1, 1), (2, 1), (3, 1), (4, 1)\}$$

...





# Representing Relations Using Matrices

## Definition

Suppose  $R$  is a relation from  $A = \{a_1, a_2, \dots, a_m\}$  to  $B = \{b_1, b_2, \dots, b_n\}$ ,  $R$  can be represented by the **matrix**  $\mathbf{M}_R = [m_{ij}]$ , where

$$m_{ij} = \begin{cases} 1 & \text{if } (a_i, b_j) \in R \\ 0 & \text{if } (a_i, b_j) \notin R \end{cases}$$

## Example

$R$  is relation from  $A = \{1, 2, 3\}$  to  $B = \{1, 2\}$ . Let  $R = \{(2, 1), (3, 1), (3, 2)\}$ , the matrix for  $R$  is

$$\mathbf{M}_R = \begin{bmatrix} 0 & 0 \\ 1 & 0 \\ 1 & 1 \end{bmatrix}$$

Determine whether the relation has certain properties (reflexive, symmetric, antisymmetric,...)



# Representing Relations Using Digraphs



## Definition

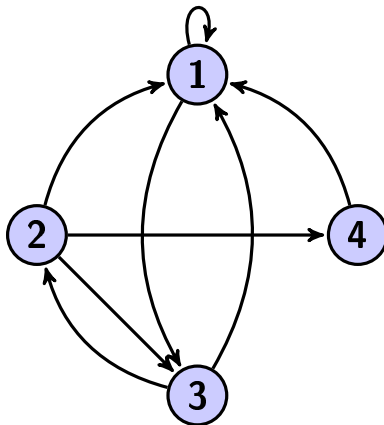
Suppose  $R$  is a relation in  $A = \{a_1, a_2, \dots, a_m\}$ ,  $R$  can be represented by the **digraph** (đồ thị có hướng)  $G = (V, E)$ , where

$$V = A$$
$$(a_i, a_j) \in E \text{ if } (a_i, a_j) \in R$$

## Example

Given a relation on  $A = \{1, 2, 3, 4\}$ ,  
 $R = \{(1, 1), (1, 3), (2, 1), (2, 3), (2, 4), (3, 1), (3, 2), (4, 1)\}$   
Draw corresponding digraph.

# Resulting digraph



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## Definition

The **closure** (*bao đóng*) of relation  $R$  with respect to **property**  $P$  is the relation  $S$  that

- i. **contains**  $R$
- ii. **has** property  $P$
- iii. is **contained in any** relation satisfying (i) and (ii).

**$S$  is the “smallest” relation satisfying (i) & (ii)**

# Reflexive Closure

## Example

Let  $R = \{(a, b), (a, c), (b, d), (d, c)\}$

The **reflexive closure** of  $R$

$\{(a, b), (a, c), (b, d), (d, c), (a, a), (b, b), (c, c), (d, d)\}$

$$R \cup \Delta$$

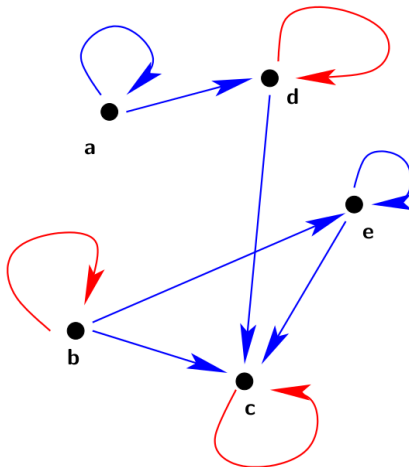
where

$$\Delta = \{(a, a) \mid a \in A\}$$

**diagonal relation** (*quan hệ đường chéo*).



# Reflexive Closure



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# Symmetric Closure



## Example

Let  $R = \{(a, b), (a, c), (b, d), (c, a), (d, e)\}$

The **symmetric closure** of  $R$

$\{(a, b), (a, c), (b, d), (c, a), (d, e), (b, a), (d, b), (e, d)\}$

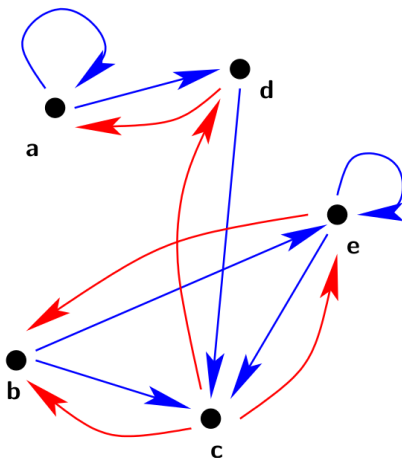
$$R \cup R^{-1}$$

where

$$R^{-1} = \{(b, a) \mid (a, b) \in R\}$$

**inverse relation** (*quan hệ ngược*).

# Symmetric Closure



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# Transitive Closure



## Example

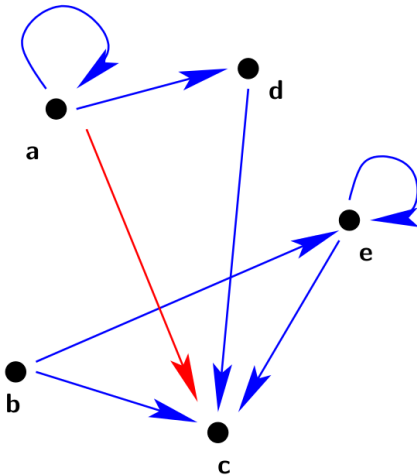
Let  $R = \{(a, b), (a, c), (b, d), (d, e)\}$

The **transitive closure** of  $R$

$\{(a, b), (a, c), (b, d), (d, e), (a, d), (b, e), (a, e)\}$

$$\bigcup_{n=1}^{\infty} R^n$$

# Transitive Closure



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# Equivalence Relations

## Definition

A relation on a set  $A$  is called an **equivalence relation** (*quan hệ tương đương*) if it is **reflexive**, **symmetric** and **transitive**.

## Example (1)

The relation  $R = \{(a, b) | a \text{ and } b \text{ are in the same provinces}\}$  is an equivalence relation.  $a$  is **equivalent** to  $b$  and vice versa, denoted  $a \sim b$ .

## Example (2)

$$R = \{(a, b) \mid a = b \vee a = -b\}$$

$R$  is an equivalence relation.

## Example (3)

$$R = \{(x, y) \mid |x - y| < 1\}$$

Is  $R$  an equivalence relation?



## Example

### Example (Congruence Modulo $m$ - Đồng dư modulo $m$ )

Let  $m$  be a positive integer with  $m > 1$ . Show that the relation

$$R = \{(a, b) \mid a \equiv b \pmod{m}\}$$

is an equivalence relation on the set of integers.

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# Equivalence Classes

## Definition

Let  $R$  be an **equivalence relation** on the set  $A$ . The set of all elements that are related to an element  $a$  of  $A$  is called the **equivalence class** (*lớp tương đương*) of  $a$ , denoted by

$$[a]_R = \{s \mid (a, s) \in R\}$$

## Example

The equivalence class of “Thủ Đức” for the equivalence relation “in the same provinces” is { “Thủ Đức”, “Gò Vấp”, “Bình Thạnh”, “Quận 10”, ... }



## Example

### Example

What are the equivalence classes of 0, 1, 2, 3 for congruence modulo 4?

### Solution:

$$[0]_4 = \{\dots, -8, -4, 0, 4, 8, \dots\}$$

$$[1]_4 = \{\dots, -7, -3, 1, 5, 9, \dots\}$$

$$[2]_4 = \{\dots, -6, -2, 2, 6, 10, \dots\}$$

$$[3]_4 = \{\dots, -5, -1, 3, 7, 11, \dots\}$$



# Equivalence Relations and Partitions



## Theorem

*Let  $R$  be an equivalence relation on a set  $A$ . These statements for elements  $a$  and  $b$  of  $A$  are equivalent:*

- i  $aRb$
- ii  $[a] = [b]$
- iii  $[a] \cap [b] \neq \emptyset$

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## Example 1

### Example

Suppose that  $S = \{1, 2, 3, 4, 5, 6\}$ . The collection of sets  $A_1 = \{1, 2, 3\}$ ,  $A_2 = \{4, 5\}$ , and  $A_3 = \{6\}$  forms a partition of  $S$ , because these sets are disjoint and their union is  $S$ .

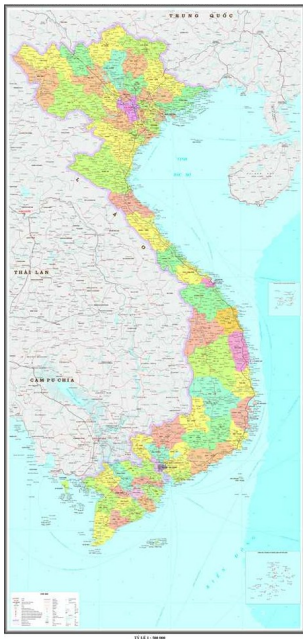
The equivalence classes of an equivalence relation  $R$  on a set  $S$  form a **partition** of  $S$ .

**Every partition** of a set can be used to form an **equivalence relation**.





## Example 2



### Example

Divides set of all cities and towns in Vietnam into set of 64 provinces. We know that:

- there are no provinces with no cities or towns
- no city is in more than one province
- every city is accounted for

### Definition

A **partition** of a Vietnam is a collection of non-overlapping non-empty subsets of Vietnam (provinces) that, together, make up all of Vietnam.

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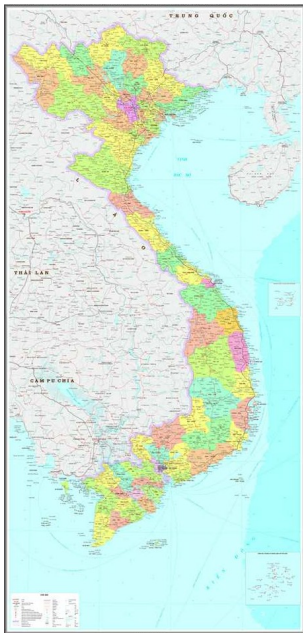
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# Relation in a Partition



- We divided based on relation

$$R = \{(a, b) | a \text{ and } b \text{ are in the same provinces}\}$$

- “Thủ Đức” is related (equivalent) to “Gò Vấp”
- “Đà Lạt” is **not** related (not equivalent) to “Long Xuyên”

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# Partial Order Relations

- Order words such that  $x$  comes before  $y$  in the dictionary
- Schedule projects such that  $x$  must be completed before  $y$
- Order set of integers, where  $x < y$

## Definition

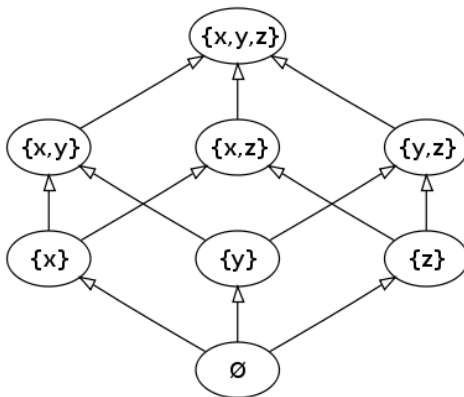
A relation  $R$  on a set  $S$  is called a **partial ordering** (có thứ tự bộ phận) if it is **reflexive**, **antisymmetric** and **transitive**. A set  $S$  together with a partial ordering  $R$  is called a partially ordered set, or **poset** (tập có thứ tự bộ phận), and is denoted by  $(S, R)$  or  $(S, \preceq)$ .

## Example

- $(\mathbb{Z}, \geq)$  is a poset
- Let  $S$  a set,  $(P(S), \subseteq)$  is a poset



## Example



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# Totally Order Relations

## Example

In the poset  $(\mathbb{Z}^+, |)$ , 3 and 9 are **comparable** (so sánh được), because  $3 \mid 9$ , but 5 and 7 are not, because  $5 \nmid 7$  and  $7 \nmid 5$ .

→ That's why we call it **partially** ordering.

## Definition

If  $(S, \preccurlyeq)$  is a poset and every two elements of  $S$  are comparable,  $S$  is called a **totally ordered** (có thứ tự toàn phần). A totally ordered set is also called a **chain** (dây xích).

## Example

The poset  $(\mathbb{Z}, \leq)$  is totally ordered.



# Maximal & Minimal Elements

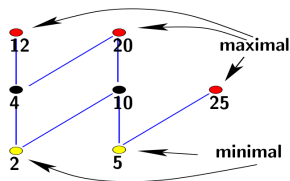


## Definition

- $a$  is **maximal** (*cực đại*) in the poset  $(S, \preceq)$  if there is no  $b \in S$  such that  $a \prec b$ .
- $a$  is **minimal** (*cực tiểu*) in the poset  $(S, \preceq)$  if there is no  $b \in S$  such that  $b \prec a$ .

## Example

Which elements of the poset  $(\{2, 4, 5, 10, 12, 20, 25\}, |)$  are minimal and maximal?



# Greatest Element & Least Element



## Definition

- $a$  is the **greatest element** (*lớn nhất*) of the poset  $(S, \preceq)$  if  $b \preceq a$  for all  $b \in S$ .
- $a$  is the **least element** (*nhỏ nhất*) of the poset  $(S, \preceq)$  if  $a \preceq b$  for all  $b \in S$ .

The greatest and least element are **unique** if it exists.

## Example

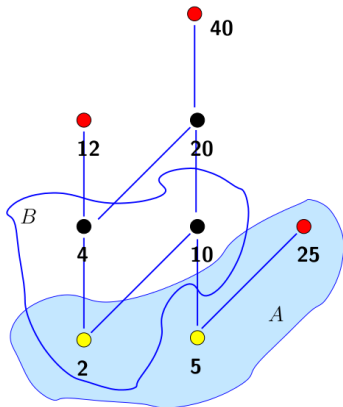
Let  $S$  be a set. In the poset  $(P(S), \subseteq)$ , the least element is  $\emptyset$  and the greatest element is  $S$ .

# Upper Bound & Lower Bound

## Definition

Let  $A \subseteq (S, \preccurlyeq)$ .

- If  $u$  is an element of  $S$  such that  $a \preccurlyeq u$  for all elements  $a \in A$ , then  $u$  is called an **upper bound** (*cận trên*) of  $A$ .
- If  $l$  is an element of  $S$  such that  $l \preccurlyeq a$  for all elements  $a \in A$ , then  $l$  is called a **lower bound** (*cận dưới*) of  $A$ .



## Example

- Subset  $A$  does **not** have upper bound and lower bound.
- The upper bound of  $B$  are 20, 40 and the lower bound is 2.

