



[Contents](#)

[Predicate Logic](#)

[Exercise](#)

# Chapter 2

## Logics (cont.)

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Nguyen An Khuong, Tran Tuan Anh, Nguyen Tien Thinh, Mai  
Xuan Toan, Tran Hong Tai  
Faculty of Computer Science and Engineering  
University of Technology - VNUHCM  
trtanh@hcmut.edu.vn

# Contents

## ① Predicate Logic

## ② Exercise

Logics (cont.)

Nguyen An Khuong,  
Tran Tuan Anh, Nguyen  
Tien Thinh, Mai Xuan  
Toan, Tran Hong Tai



Contents

Predicate Logic

Exercise

# Course outcomes



## Contents

Predicate Logic

Exercise

### Course learning outcomes

L.O.1	Understanding of logic and discrete structures
	<a href="#">L.O.1.1 – Describe definition of propositional and predicate logic</a>
	L.O.1.2 – Define basic discrete structures: set, mapping, graphs
L.O.2	Represent and model practical problems with discrete structures
	<a href="#">L.O.2.1 – Logically describe some problems arising in Computing</a>
	L.O.2.2 – Use proving methods: direct, contrapositive, induction
	L.O.2.3 – Explain problem modeling using discrete structures
L.O.3	Understanding of basic probability and random variables
	L.O.3.1 – Define basic probability theory
	L.O.3.2 – Explain discrete random variables
L.O.4	Compute quantities of discrete structures and probabilities
	L.O.4.1 – Operate (compute/ optimize) on discrete structures
	L.O.4.2 – Compute probabilities of various events, conditional ones, Bayes theorem

# Limits of Propositional Logic

- $x > 3$
- All square numbers are not prime numbers. 100 is a square number. Therefore 100 is not a prime number.

Logics (cont.)

Nguyen An Khuong,  
Tran Tuan Anh, Nguyen  
Tien Thinh, Mai Xuan  
Toan, Tran Hong Tai



Contents

Predicate Logic

Exercise



## Definition

A predicate (*vị từ*) is a statement **containing** one or more **variables**. If **values are assigned** to all the variables in a predicate, the resulting statement is a **proposition** (*mệnh đề*).

- $x > 3 \rightarrow P(x)$
- $5 > 3 \rightarrow P(5)$
- A predicate with  $n$  variables  $P(x_1, x_2, \dots, x_n)$

Example:

- $x > 3$  (**predicate**)
- $5 > 3$  (**proposition**)
- $2 > 3$  (**proposition**)

# Truth value

- $x > 3$  is true or false?
- $5 > 3$
- For every number  $x$ ,  $x > 3$  holds
- There is a number  $x$  such that  $x > 3$

Logics (cont.)

Nguyen An Khuong,  
Tran Tuan Anh, Nguyen  
Tien Thinh, Mai Xuan  
Toan, Tran Hong Tai



Contents

Predicate Logic

Exercise

# Quantifiers

- $\forall$ : Universal – *Với mọi*
  - $\forall x P(x) = P(x)$  is T for all  $x$
- $\exists$ : Existential – *Tồn tại*
  - $\exists x P(x) =$  There exists an element  $x$  such that  $P(x)$  is T
- We need a **domain of discourse** for variable





### Example

Let  $P(x)$  be the statement “ $x < 2$ ”. What is the truth value of the quantification  $\forall xP(x)$ , where the domain consists of all real number?

- $P(3) = 3 < 2$  is false
- $\Rightarrow \forall xP(x)$  is false
- 3 is a **counterexample** (*phản ví dụ*) of  $\forall xP(x)$

### Example

What is the truth value of the quantification  $\exists xP(x)$ , where the domain consists of all real number?





## Example

Express the statement "Some student in this class comes from Central Vietnam."

### Solution 1

- $M(x) = x$  comes from Central Vietnam
- Domain for  $x$  is **the students in the class**
- $\exists x M(x)$

### Solution 2

- Domain for  $x$  is all people
- ...

# Negation of Quantifiers

Statement	Negation	Equivalent form
$\forall x P(x)$	$\neg(\forall x P(x))$	$\exists x \neg P(x)$
$\exists x P(x)$	$\neg(\exists x P(x))$	$\forall x \neg P(x)$

## Example

- **All** CSE students study Discrete Math 1
- Let  $C(x)$  denote "x is a CSE student"
- Let  $S(x)$  denote "x studies Discrete Math 1"
- $\forall x : C(x) \rightarrow S(x)$
- $\exists x : \neg(C(x) \rightarrow S(x)) \equiv \exists x : C(x) \wedge \neg S(x)$
- **There is** a CSE student who does not study Discrete Math 1.



## Another Example

### Example

Translate these:

- All lions are fierce.
- Some lions do not drink coffee.
- Some fierce creatures do not drink coffee.

### Solution

Let  $P(x)$ ,  $Q(x)$  and  $R(x)$  be the statements “ $x$  is a lion”, “ $x$  is fierce” and “ $x$  drinks coffee”, respectively.

- $\forall x(P(x) \rightarrow Q(x))$ .
- $\exists x(P(x) \wedge \neg R(x))$ .
- $\exists x(Q(x) \wedge \neg R(x))$ .



# The Order of Quantifiers

- The **order** of quantifiers is **important**, unless **all** the quantifiers are universal quantifiers or all are existential quantifiers
- Read from left to right, apply from inner to outer

## Example

$$\forall x \forall y (x + y = y + x)$$

**T** for all  $x, y \in \mathbb{R}$

## Example

$$\forall x \exists y (x + y = 0) \text{ is } \mathbf{T},$$

while

$$\exists y \forall x (x + y = 0) \text{ is } \mathbf{F}$$



# Translating Nested Quantifiers

Logics (cont.)

Nguyen An Khuong,  
Tran Tuan Anh, Nguyen  
Tien Thinh, Mai Xuan  
Toan, Tran Hong Tai



Contents

Predicate Logic

Exercise

## Example

$$\forall x (C(x) \vee \exists y (C(y) \wedge F(x, y)))$$

Provided that:

- $C(x)$ :  $x$  has a computer,
- $F(x, y)$ :  $x$  and  $y$  are friends,
- $x, y \in$  all students in your school.

## Answer

For every student  $x$  in your school,  $x$  has a computer or there is a student  $y$  such that  $y$  has a computer and  $x$  and  $y$  are friends.

# Translating Nested Quantifiers

Logics (cont.)

Nguyen An Khuong,  
Tran Tuan Anh, Nguyen  
Tien Thinh, Mai Xuan  
Toan, Tran Hong Tai



Contents

Predicate Logic

Exercise

## Example

$\exists x \forall y \forall z \quad (((F(x, y) \wedge F(x, z) \wedge (y \neq z)) \rightarrow \neg F(y, z)))$

Provided that:

- $F(x, y)$ :  $x, y$  are friends
- $x, y, z \in$  all students in your school.

## Answer

There is a student  $x$ , so that for every student  $y$ , every student  $z$  not the same as  $y$ , if  $x$  and  $y$  are friends, and  $x$  and  $z$  are friends, then  $y$  and  $z$  are not friends.

# Translating into Logical Expressions

Logics (cont.)

Nguyen An Khuong,  
Tran Tuan Anh, Nguyen  
Tien Thinh, Mai Xuan  
Toan, Tran Hong Tai



Contents

Predicate Logic

Exercise

## Example

- 1 “**There is** a student in the class has visited Hanoi”.
- 2 “**Every** student in the class has visited Nha Trang **or** Vung Tau”.

## Answer

Assume:

$C(x)$  :  $x$  has visited Hanoi

$D(x)$  :  $x$  has visited Nha Trang

$E(x)$  :  $x$  has visited Vung Tau

We have:

- 1  $\exists x C(x)$
- 2  $\forall x (D(x) \vee E(x))$

# Translating into Logical Expressions

Logics (cont.)

Nguyen An Khuong,  
Tran Tuan Anh, Nguyen  
Tien Thinh, Mai Xuan  
Toan, Tran Hong Tai



[Contents](#)

[Predicate Logic](#)

[Exercise](#)

## Example

If a person is a woman and a parent, then this person is mother of someone.

## Solution

We define:

- $W(x)$  :  $x$  is woman
- $P(x)$  :  $x$  is a parent
- $M(x, y)$  :  $x$  is mother of  $y$

We have:  $\forall x((W(x) \wedge P(x)) \rightarrow \exists y M(x, y))$

## Example

"Every people has only one best friend."

Assume:

- $B(x, y)$  :  $y$  is the best friend of  $x$



# Translating into Logical Expressions

Logics (cont.)

Nguyen An Khuong,  
Tran Tuan Anh, Nguyen  
Tien Thinh, Mai Xuan  
Toan, Tran Hong Tai



[Contents](#)

[Predicate Logic](#)

[Exercise](#)

## Example

"Every people has only one best friend."

Assume:

- $B(x, y) : y$  is the best friend of  $x$

## Solution

$$\forall x \exists y \forall z (B(x, y) \wedge ((y \neq z) \rightarrow \neg B(x, z)))$$



## Example

- If I have a girlfriend, I will take her to go shopping.
- Whenever I and my girlfriend go shopping and that day is a special day, I will surely buy her some expensive gift.
- If I buy my girlfriend expensive gifts, I will eat noodles for a week.
- Today is March 8.
- March 8 is such a special day.
- Therefore, if I have a girlfriend,...
- I will eat noodles for a week.

# Propositional Rules of Inferences

Rule of Inference	Name
$\begin{array}{l} p \\ p \rightarrow q \\ \hline \therefore q \end{array}$	Modus ponens
$\begin{array}{l} \neg q \\ p \rightarrow q \\ \hline \therefore \neg p \end{array}$	Modus tollens
$\begin{array}{l} p \rightarrow q \\ q \rightarrow r \\ \hline \therefore p \rightarrow r \end{array}$	Hypothetical syllogism ( <i>Tam đoạn luận giả định</i> )
$\begin{array}{l} p \vee q \\ \neg p \\ \hline \therefore q \end{array}$	Disjunctive syllogism ( <i>Tam đoạn luận tuyển</i> )

Logics (cont.)

Nguyen An Khuong,  
Tran Tuan Anh, Nguyen  
Tien Thinh, Mai Xuan  
Toan, Tran Hong Tai



Contents

Predicate Logic

Exercise

# Propositional Rules of Inferences

Rule of Inference	Name
$\frac{p}{\therefore p \vee q}$	Addition ( <i>Quy tắc cộng</i> )
$\frac{p \wedge q}{\therefore p}$	Simplification ( <i>Rút gọn</i> )
$\frac{p}{q}$ $\frac{q}{\therefore p \wedge q}$	Conjunction ( <i>Kết hợp</i> )
$\frac{p \vee q}{\neg p \vee r}$ $\frac{\neg p \vee r}{\therefore q \vee r}$	Resolution ( <i>Phân giải</i> )

Logics (cont.)

Nguyen An Khuong,  
Tran Tuan Anh, Nguyen  
Tien Thinh, Mai Xuan  
Toan, Tran Hong Tai



[Contents](#)

[Predicate Logic](#)

[Exercise](#)



## Example

If it rains today, then we will not have a barbecue today. If we do not have a barbecue today, then we will have a barbecue tomorrow. Therefore, if it rains today, then we will have a barbecue tomorrow.

## Solution

- $p$ : It is raining today
- $q$ : We will not have a barbecue today
- $r$ : We will have barbecue tomorrow

$$p \rightarrow q$$

$$q \rightarrow r$$

$$\hline \therefore p \rightarrow r$$

Hypothetical syllogism

[Contents](#)[Predicate Logic](#)[Exercise](#)

## Example

- It is not sunny this afternoon ( $\neg p$ ) and it is colder than yesterday ( $q$ )
- We will go swimming ( $r$ ) only if it is sunny
- If we do not go swimming, then we will take a canoe trip ( $s$ )
- If we take a canoe trip, then we will be home by sunset ( $t$ )
- We will be home by sunset ( $t$ )

1.  $\neg p \wedge q$  Hypothesis
2.  $\neg p$  Simplification using (1)
3.  $r \rightarrow p$  Hypothesis
4.  $\neg r$  Modus tollens using (2) and (3)
5.  $\neg r \rightarrow s$  Hypothesis
6.  $s$  Modus ponens using (4) and (5)
7.  $s \rightarrow t$  Hypothesis
8.  $t$  Modus ponens using (6) and (7)



## Definition

Fallacies (*ngụy biện*) resemble rules of inference but are based on contingencies rather than tautologies.

## Example

If you do correctly every questions in mid-term exam, you will get 10 grade. You got 10 grade.

Therefore, you did correctly every questions in mid-term exam.

Is  $[(p \rightarrow q) \wedge q] \rightarrow p$  a tautology?

# Rules of Inference for Quantified Statements

Logics (cont.)

Nguyen An Khuong,  
Tran Tuan Anh, Nguyen  
Tien Thinh, Mai Xuan  
Toan, Tran Hong Tai



[Contents](#)

[Predicate Logic](#)

[Exercise](#)

Rule of Inference	Name
$\frac{\forall x P(x)}{\therefore P(c)}$	Universal instantiation ( <i>Cụ thể hóa phổ quát</i> )
$\frac{P(c) \text{ for an arbitrary } c}{\therefore \forall x P(x)}$	Universal generalization ( <i>Tổng quát hóa phổ quát</i> )
$\frac{\exists x P(x)}{\therefore P(c) \text{ for some element } c}$	Existential instantiation ( <i>Cụ thể hóa tồn tại</i> )
$\frac{P(c) \text{ for some element } c}{\therefore \exists x P(x)}$	Existential generalization ( <i>Tổng quát hóa tồn tại</i> )



[Contents](#)[Predicate Logic](#)[Exercise](#)

## Example

- A student in this class has not gone to class
- Everyone in this class passed the first exam
- Someone who passed the first exam has not gone to class

## Hint

- $C(x)$ :  $x$  is in this class
- $B(x)$ :  $x$  has gone to class
- $P(x)$ :  $x$  passed the first exam
- Premises???

[Contents](#)[Predicate Logic](#)[Exercise](#)

- |                                       |                                     |
|---------------------------------------|-------------------------------------|
| 1. $\exists x(C(x) \wedge \neg B(x))$ | Premise                             |
| 2. $C(a) \wedge \neg B(a)$            | Existential instantiation from (1)  |
| 3. $C(a)$                             | Simplification from (2)             |
| 4. $\forall x(C(x) \rightarrow P(x))$ | Premise                             |
| 5. $C(a) \rightarrow P(a)$            | Universal instantiation from (4)    |
| 6. $P(a)$                             | Modus ponens from (3) and (5)       |
| 7. $\neg B(a)$                        | Simplification from (2)             |
| 8. $P(a) \wedge \neg B(a)$            | Conjunction from (6) and (7)        |
| 9. $\exists x(P(x) \wedge \neg B(x))$ | Existential generalization from (8) |



Given the predicate  $p(x) : "x^2 - 3x + 2 = 0"$ . What is the truth value (chân trị) of the following propositions:

- a)  $p(0)$
- b)  $p(1)$
- c)  $p(2)$
- d)  $\exists x, p(x)$
- e)  $\forall x, p(x)$

[Contents](#)[Predicate Logic](#)[Exercise](#)

Let  $x, y \in \mathbb{Z}^+$ , and the predicate:  $p(x, y)$ : " $x$  is a divisor of  $y$ "  
Determine the truth value of the following propositions:

- a)  $p(2, 3)$
- b)  $p(2, 6)$
- c)  $\forall y, p(1, y)$
- d)  $\forall x, p(x, x)$
- e)  $\forall x \exists y, p(x, y)$
- f)  $\exists y \forall x, p(x, y)$
- g)  $\forall x \forall y, (p(x, y) \wedge p(y, x)) \rightarrow (x = y)$
- h)  $\forall x \forall y \forall z (p(x, y) \wedge p(y, z)) \rightarrow (p(x, z))$



Provided that:

- $F(x, y)$  :  $x$  is father of  $y$ ,
- $M(x, y)$  :  $x$  is mother of  $y$ ,
- $S(x, y)$  :  $x$  is sister of  $y$ ,
- $B(x, y)$  :  $x$  is brother of  $y$ ,
- $H(x, y)$  :  $x$  is spouse (wife/husband) of  $y$ ,
- $O(x, y)$  :  $x$  is elder than  $y$ .

**Express each of these statements using predicates:**

- 'He (a person) has an elder sister and younger brother'.
- 'All of her brothers are younger than her'.
- 'Thuyen has only one husband' (Thuyen is a private name).
- 'One of his sisters is younger than him'.
- 'Everyone has grandfather, grandmother, maternal grandfather, maternal grandmother'.
- 'A father of a person cannot be a mother of other ones'.



## Solutions:

- a) 'He (a person) has an elder sister and younger brother'.  
 $\exists x \exists y (S(x, m) \wedge O(x, m) \wedge B(y, m) \wedge \neg O(y, m)).$
- b) 'All of her brothers are younger than her'.  
 $\forall x (B(x, m) \rightarrow \neg O(x, m)).$
- c) 'Thuyen has only one husband' (Thuyen is a private name).  
 $\exists x \forall y (H(x, \text{Thuyen}) \wedge H(y, \text{Thuyen}) \rightarrow (x = y))$   
or  $\exists x \forall y (H(x, \text{Thuyen}) \wedge (x \neq y) \rightarrow \neg H(y, \text{Thuyen})).$
- d) 'One of his sisters is younger than him'.  
 $\exists x \forall y (S(x, m) \wedge \neg O(x, m) \wedge S(y, m) \wedge (x \neq y) \rightarrow O(y, m)).$
- e) 'Everyone has grandfather, grandmother, maternal grandfather, maternal grandmother'.  $\forall x \exists y \exists z \exists y_1 \exists y_2 \exists z_1 \exists z_2$   
 $(F(y, x) \wedge M(z, x) \wedge F(y_1, y) \wedge M(y_2, y) \wedge F(z_1, z) \wedge M(z_2, z)).$
- f) 'A father of a person cannot be a mother of other ones'.  
 $\exists x \exists y \forall z (F(x, y) \rightarrow \neg M(x, z)).$

[Contents](#)[Predicate Logic](#)[Exercise](#)

### Translating the following nested quantifiers:

- a)  $B(c, m) \wedge (O(c, m) \vee O(m, c)).$
- b)  $B(c, m) \wedge F(a, m) \rightarrow O(a, c) \wedge F(a, c).$
- c)  $\forall x \forall y (S(x, m) \wedge B(c, y) \rightarrow x = y).$
- d)  $\exists x ((S(x, m) \vee H(c, x)) \vee \exists x (H(x, m) \wedge O(x, m))).$
- e)  $\forall x \forall y (S(x, m) \wedge S(y, m) \rightarrow O(x, y) \vee O(y, x))$



Given a predicate  $N(x)$  " $x$  has been to Da Lat" with the domain is the all students in Mathematics class. Translate the following predicates into English

- a)  $\exists x N(x)$
- b)  $\forall x N(x)$
- c)  $\neg \exists x N(x)$
- d)  $\exists x \neg N(x)$
- e)  $\neg \forall x N(x)$
- f)  $\forall x \neg N(x)$

- a) There is a student in this class has been to Da Lat.
- b) All students in Math class have been to Da Lat.
- c) There is no exists a student in Math class has gone to Da Lat.
- d) There is a student in this class has never gone to Da Lat.
- e) Not all students in Math class have ever been to Da Lat.
- f) All students in Math class have never been to Da Lat.





Given the predicate  $N(x)$  " $x$  studies more than 5 hours in class every weekday" with the domain is the all students in Mathematics class. Express the following predicates:

- a)  $\exists x N(x)$
- b)  $\forall x N(x)$
- c)  $\exists x \neg N(x)$
- d)  $\forall x \neg N(x)$

[Contents](#)[Predicate Logic](#)[Exercise](#)

What is the propositional formula for the following pseudo code:

```
for (i = 0; i < numObjects; i++) {  
    Object x = Objects(i);  
    if isMushroom(x)  
        if isPoisonous(x) && isPurple(x)  
            return false;  
}  
return true;
```

- There are no mushrooms that are poisonous and purple.
- $\forall x \text{Mushroom}(x) \rightarrow \neg(\text{Poisonous}(x) \wedge \text{Purple}(x))$

[Contents](#)[Predicate Logic](#)[Exercise](#)

What is the propositional formula for the following pseudo code:

```
for (i=0; i<numObjects; i++) {  
    Object x = Objects(i);  
    if isMushroom(x) && isPoisonous(x) && isPurple(x)  
        return true;  
}  
return false;
```

- There is a mushroom that is purple and poisonous.
- $\exists x \text{Mushroom}(x) \wedge \text{Poisonous}(x) \wedge \text{Purple}(x)$



Giving the following pseudo code:

```
//— Look for first match
for (x=0; x<numKids; x++)
  if isParent(Peter, kids[x])
    match1Found = true;
```

$$\exists x \exists y ( \text{isParent}(P, x) \wedge \neg (x = y) )$$

```
//— Now look for a second match
for (y=0; (y<numKids)&&(y!=x); y++)
  if isParent(Peter, kids[y])
    match2Found = true;
```

```
return match1Found && match2Found;
```

Knowing that: kids array has 3 elements: { Alice, Bob, Charles }  
 and Peter only have 1 child Alice.

What is the propositional formula for "Peter has at least 2 children".

$$\exists x \exists y (\text{ParentOf}(\text{Peter}, x) \wedge \text{ParentOf}(\text{Peter}, y) \wedge \neg(x = y))$$



Let  $P(x)$  be " $x$  can speak Russian" and  $Q(x)$  be " $x$  can use Java".  
Formalize the following:

Giving the space is all students in your university.

- a) There is a student in your university that can speak Russian and can use Java.
- b) There is a student in your university that can speak Russian but can't use Java.
- c) Every student in your university can speak Russian or can use Java.
- d) None of the student in your university can speak Russian or can use Java.

- a)  $\exists x(P(x) \wedge Q(x))$
- b)  $\exists x(P(x) \wedge \neg Q(x))$
- c)  $\forall x(P(x) \vee Q(x))$
- d)  $\forall x\neg(P(x) \wedge Q(x))$



Let  $L(x,y)$  be "x love y", where the space of  $x$  and  $y$  is the set of all people in the world. Use logical quantifier to express the following:

- a) Everybody loves Jerry.
- b) Everybody loves someone.
- c) There is a person who everybody loves.
- d) Nobody loves everybody.
- e) There is someone Lydica doesn't love.
- f) There is someone nobody loves.
- g) There is exact one person everybody loves.
- h) There are exact two person Lynn loves.
- i) Everybody loves themselves.
- j) There is a person who love nobody but himself.



Giving the following:

- .  $\neg P(x)$ : "x is a math problem".
- .  $\neg Q(x)$ : "x is hard" (based on a well-defined standard).
- .  $\neg R(x)$ : "x is easy" (based on a well-defined standard - same as above).
- .  $\neg S(x)$ : "x is not solvable".

Translate the following propositional formulas to natural English

- a)  $\forall x(P(x) \rightarrow (Q(x) \iff \neg R(x)))$
- b)  $\exists x(S(x) \wedge \neg P(x))$

There are many ways to translate a formula to natural language and the following is one of them

- a) If x is a math problem, to say x is hard is the same as saying x is not easy.
- b) There is unsolvable non-math problem.

[Contents](#)[Predicate Logic](#)[Exercise](#)

Translate the following propositional formulas to natural English where:

$F(p)$  is “Printer  $p$  is broken”,

$B(p)$  is “Printer  $p$  is currently printing another document”,

$L(j)$  is “Printing job  $j$  is lost”,

and  $Q(j)$  is “Printing job  $j$  is in queue.”

a)  $\exists p(F(p) \wedge B(p)) \rightarrow \exists jL(j)$

b)  $\forall pB(p) \rightarrow \exists jQ(j)$

c)  $\exists j(Q(j) \wedge L(j)) \rightarrow \exists pF(p)$

d)  $(\forall pB(p) \wedge \forall jQ(j)) \rightarrow \exists jL(j)$



[Contents](#)[Predicate Logic](#)[Exercise](#)

Formalize the following sentences:

- a) Nobody is perfect.
- b) not everyone is perfect.
- c) All your friends are perfect.
- d) At least one of your friend is perfect.
- e) Everybody is your friend and they are perfect.
- f) Not everybody is your friend or there is somebody not perfect.

Giving:  $C(x)$ :  $x$  is perfect.

$D(x)$ :  $x$  is your friend.

$E(x)$ :  $x$  is someone else.



Giving the following Predicate:

- $P(x)$ : Program  $x$  satisfies ABET standard.
- $Q(x,y)$ : Program  $x$  has the same educational goal as program  $y$ .
- $R(x)$ : Educational outcome from program  $x$  is verifiable.

Which of the following formalize this sentence : "Every program that has the same educational goal as a ABET satisfied program and verifiable Educational outcome also satisfies ABET standard"

- A  $\forall x(P(x) \wedge \neg Q(x)) \rightarrow \exists x(R(x))$
- B  $\forall x(\exists y(Q(x,y) \wedge P(y) \wedge R(x)) \rightarrow P(x))$
- C  $\forall x(\exists y(Q(x,y) \wedge P(y) \wedge R(x)) \rightarrow P(x) \vee R(x))$
- D  $\forall x(\forall y(Q(x,y) \wedge P(y) \vee R(x)) \rightarrow P(x))$



Let:

- $P(x, y)$ :  $x$  is parent of  $y$ .
- $M(x)$ :  $x$  is male .

Given:

$F(v, w) = M(v) \wedge \exists x \exists y (P(x, y) \wedge P(x, v) \wedge (y \neq v) \wedge P(y, w))$ ,  
then  $F(v, w)$  means:

- ☐ A)  $v$  is brother of  $w$
- ☐ B)  $v$  is cousin of  $w$
- ☒ C)  $v$  is uncle of  $w$
- ☐ D)  $v$  is grand father of  $w$

[Contents](#)[Predicate Logic](#)[Exercise](#)

Formalize the following sentences using predicate logic:

- a) When a hard drive has less than 30GB free space, a warning will be issued to all the users.
- b) Do not back up the files if anyone is logging in the system.
- c) YouTube's videos will be buffered if there are at least 8MB memory and 56kb/s line rate.
- d) Few computer student is good at programming.
- e) No computer student is not hard working.
- f) Not all computer students are smart.
- g) All the Pompeians are either loyal to or hate Caesar.
- h) Everyone is loyal to someone.
- i) People only want to assassinate the dictator whom they are not loyal to.