

# COMP1002

# DATA STRUCTURES AND

# ALGORITHMS

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## LECTURE 3: STACKS, QUEUES AND OBJECTS

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# Objectives

- Revise Object orientation
- Introduce Abstract Data Types
- Provide first examples of ADTs – Stacks and Queues
- Discuss applications of Stacks/Queues incl. Equation Solving
  - Postfix evaluation
  - Infix to postfix conversion

# OBJECT ORIENTATION

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Revision slides from OOPD/FOP

# Object-Orientation

- In object-oriented programming, we bundle the behaviour (methods) and data (attributes) together
- Benefits:
  - OO protects data from being used incorrectly
  - Increases code reuse (fewer errors)
  - Makes code easier to read and maintain
  - Objects "know" how to respond to requests
  - Relates to how objects function in the real world

# Classes – Specifying Objects

- Before we can use an object, we need to describe it as a **class** (of objects).
  - Similar to how we define a function once and use it multiple times
- The class specifies the state and behaviour an object can have:
  - State: what the object is
    - attributes or member fields
  - Behaviour: what the object does
    - methods or functions

# Encapsulation

- A (an object of a) class makes use of the "information hiding" principle
  - Communication with the rest of the software system is clearly defined
    - methods are the means for communication
  - Its obligations to the software system are clearly defined
    - what services the class offers (via data and methods)
  - Implementation details should be hidden from the user
    - don't need to know how it does things to use it

# Class Specification

- Must include:
  - Details of the communication with the rest of the software system (method names)
  - The exact data representation required
  - Exactly how the required functionality is to be achieved (method implementation)



# Classes and Objects

- An object is an **instance of** a class
- The class definition provides a template for an object
- An object gives details for a particular instance

Specific cat = instance  
"Oogie" of class "cat"

Generic cat = class "cat"



# Class roles

- Every class is designed with a specific **role** in mind.
- The total set of functional requirements for a software system is broken down into a set of tasks
- Collections of tasks are grouped together and mapped to roles
- Roles are mapped to specific classes

# Class Responsibility

- Take the requirements for a software application:
  - Identify the classes required
  - Assign specific Responsibilities to each class
  - Determine relationships between classes (see later)
  - Repeat the above steps until the design is correct
- Each responsibility should be handled by that class and no other
  - Example: If a responsibility for keeping track of a person's name is assigned to a class called PersonClass then:
    - No other class should have this information
    - Other classes which need this information should refer to this class when the information is required

# Comparison to non-OO design

- In a top-down procedural approach, we design an algorithm by starting with a main module and using step-wise refinement to determine the processing steps
- Some of these steps get refined into sub modules and the process repeats until the design is refined enough to code
- Under Object Orientation this all changes...

# OO design

- Before the algorithm is designed:
  - The classes are identified
  - Each class is assigned role(s) or responsibilities
  - The required sub modules are designed (i.e. Constructors, accessors, etc)
  - Each Class is thoroughly tested via a test harness
- Finally, the main algorithm and any required sub modules is designed (making use of the developed classes in the process)

# Nouns and Verbs

- Like algorithm design, the determination of classes is still a bit of an art form
- One simple technique is the nouns and verb approach:
  - Nouns are mapped to classes
  - Verbs are mapped to sub modules within classes
  - The definition of noun and verb gets stretched to cover collections of words
  - Result is that:
    - Sub module names should always describe an action (i.e. getName)
    - Class names should always describe a thing (e.g. PersonClass)
- It is important to note that the set of classes proposed will change over the design phase

# Object Communication

- Sometimes referred to as **message passing**:
  - When an object of one class calls an object of another class it is passing a message (i.e. A request to the object to perform some task)
- The [public] methods must provide the functionality required for the class to fulfill its role.
- There are five categories of methods in a class:
  - The Constructors
  - The Accessor Methods (aka Interrogative Methods)
  - The Mutator Methods (aka Informative Methods)
  - Doing Methods (aka Imperative Methods)
  - [Private] methods

# Classes in Python

- Order your code consistently
  - Declare the components of each class in the following order:
    - Declarations for class constants and variables (global to the class)
    - Declarations for the Constructors (`__init__`)
    - Declarations of **instance variables** (local to each instance, usually in `_init_`)
      - e.g. `self.myVar = value`
    - *Accessor methods*
    - *Mutator methods*
    - Doing methods ("public")
    - Internal methods ("private")
- } Python instance and class variables are public, so basic set/gets are not req'd*



# Classes in Python

- Note that everything in Python is "public" (unlike Java, C++) so we can only **treat** methods and data as private
- Use `_methodName` to indicate "private methods"
- Put the class files in a separate python file, e.g. `DSAShield.py`
- Your programs will then import from `DSAShield` as needed
- Unit tests (testing you classes/methods)
  - Option 1: Separate `UnitTestsDSAShield.py`
  - Option 2: Include tests in `DSAShield.py` using

```
if __name__ == "__main__":  
    <tests in here>
```

# Example: song

```
class Song():
```

```
    def __init__(self, lyrics):  
        self.lyrics = lyrics
```

```
    def sing_me_a_song(self):  
        for line in self.lyrics:  
            print(line)
```

```
lumberjack = Song(["I'm a lumberjack and I'm OK",  
                  "I sleep all night",  
                  "And I work all day"])
```

```
spam = Song(["SPAM, SPAM, SPAM, SPAM",  
            "spam, spam, spam, spam"])
```

```
lumberjack.sing_me_a_song()  
spam.sing_me_a_song()
```

Instance  
variable

Object of  
class Song

Song: lumberjack

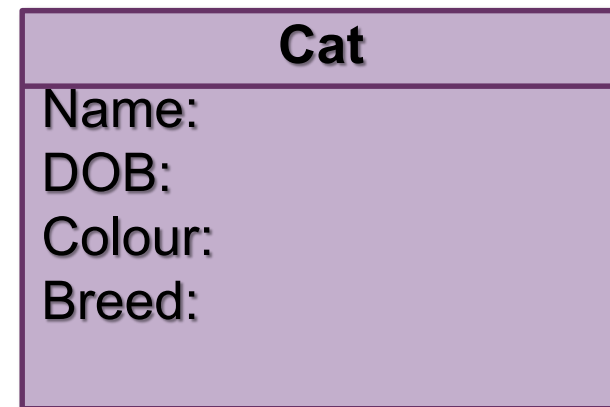
lyrics: ["I'm a lumberjack and  
I'm OK", "I sleep all night",  
"And I work all day"]

# Self

- Why do I need self when I make `__init__` or other functions for classes?
- If you don't have self, then code like `cheese = 'Gorgonzola'` is ambiguous.
- That code isn't clear about whether you mean the *instance's* cheese attribute/variable, *or* a local variable named cheese.
- With `self.cheese = 'Gorgonzola'` it's very clear you mean the instance attribute `self.cheese`.
- You can use any variable name, but self is the convention.

# OO Design...Where to begin?

- Find your objects
- If we wanted to keep track of our household animals: cats, dogs and birds
- We could make classes for cats, dogs and birds
- For each animal, we might track:
  - name
  - date of birth
  - colour
  - breed



# Test our objects out...

## CAT

Name: Oogie  
DOB: 1/1/2006  
Colour: Grey  
Breed: Fluffy



## DOG

Name: Dude  
DOB: 1/1/2011  
Colour: Brown  
Breed: Jack Russell



## BIRD

Name: Big Bird  
DOB: 10/11/1969  
Colour: Yellow  
Breed: Canary

# CLASS RELATIONSHIPS

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# Goals of Object-Orientation

- Reuse / Extensibility
  - Reuse: each class provides its functionality to other classes
  - Can inherit from a class to reuse/extend its functionality
- Modularization - low coupling, high cohesion
  - Objects should be responsible for their own data state
  - Objects should represent a single concept and all methods should relate to that concept (high cohesion)
  - Only the object's interface should matter to a user of that object, not the details of its implementation (low coupling)
- **Note:** *many of these slides are from Object-Oriented Program Design*

# Class Relationships

- The classes of objects which communicate with each other via message passing share some form of relationship (association):
  - Aggregation
  - Composition
  - Inheritance
  - Other



# Class Relationships

- Aggregation:
  - One class is declared as a class field within the other class
  - Communication is one way (most of the time?), from class to class field
- Composition:
  - One class is included as part of the other class
  - The included class does not exist without the host class

# Class Relationships

- Inheritance:
  - One class is a descendant of another class
  - Uses polymorphism, method overloading or direct references to the superclass to communicate.
  - Communication is one way, from child to parent (sound familiar!!)
- Other:
  - Where objects of one class are related to another in a manner which is NOT aggregation or inheritance.
  - These other relationships will be discussed in future units.

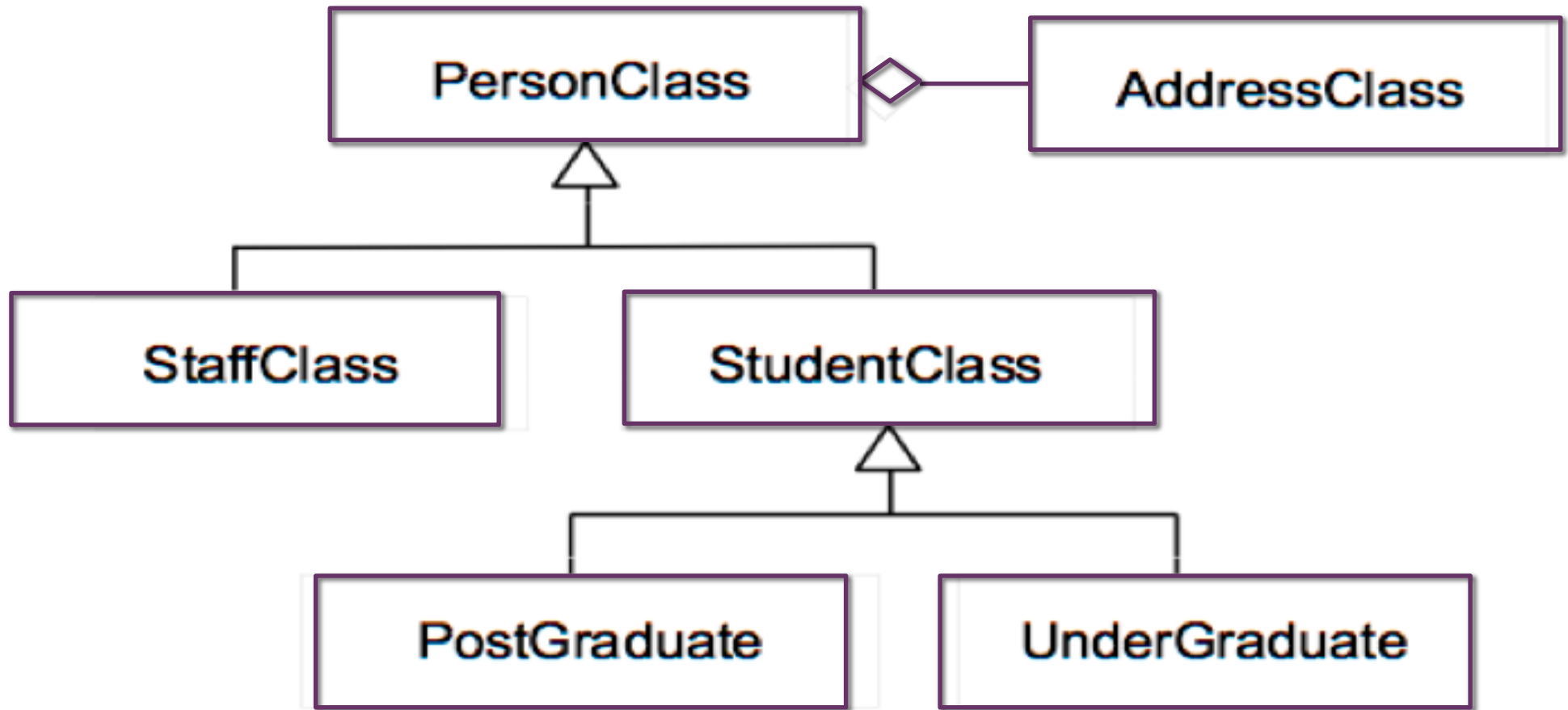
# Object Communication

- Also referred to as message passing:
- When an object of one class calls a method in an object of another class it is passing a message
- A request to the object to perform some task

# Modelling Languages

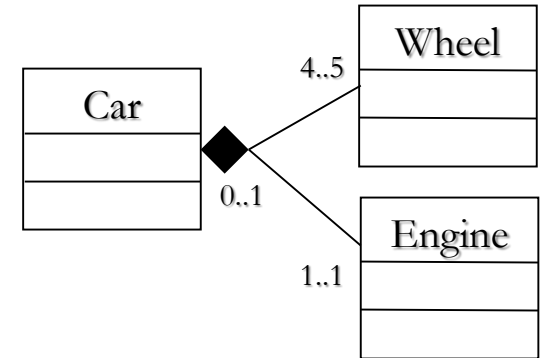
- Used to show the relationships between different classes and different instances of classes (i.e. objects) in a particular software
- Usually graphical
- Most commonly adopted methodology is known as **UML**:
  - **Unified**: a union of the approaches put forward by Grady Booch, James Rumbaugh and Ivar Jacobson
  - **Modelling**: a graphical representation (or model) of an OO software design
  - **Language**: provides a standard way of expressing object relationships (i.e. contains rules for syntax & semantics)
- Software Engineering units teach UML and OO software design.
- For now we will simply look at the UML notation for class diagrams - describing inheritance and aggregation/composition.

# Uni People Example

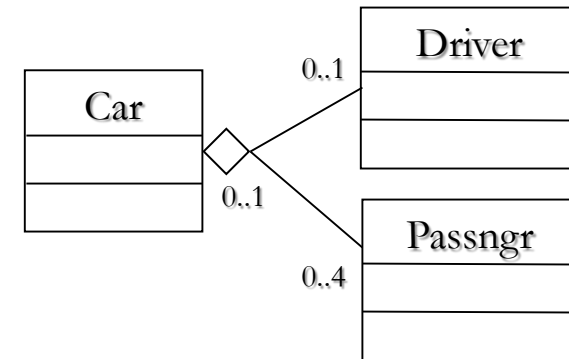


# Class Relationships (1)

- Composition
  - “**has-a**” or “whole-part” relationship
    - UML: Shown with solid diamond beside container class
    - e.g., Car “has-a” Wheel
  - Strong lifecycle dependency between classes
    - Car is not a car without four Wheels and an Engine
    - When Car is destroyed, so are the Wheels and Engine
  - In code:
    - Car would have Wheel and Engine as class fields

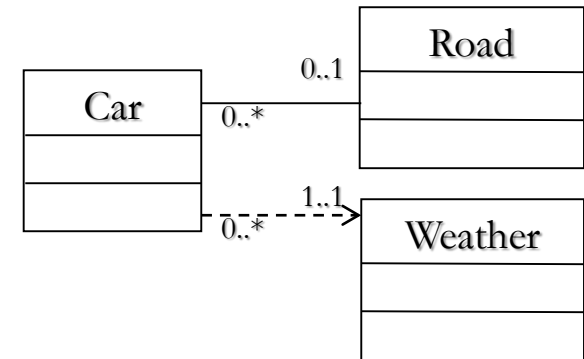


# Class Relationships (2)



- Aggregation
  - Weaker form of composition, but is still “**has-a**”
    - UML: Shown with open/unfilled diamond beside container
  - Lifecycle dependency usually not strong
    - Car does not always have a driver
    - When Car is destroyed driver and passengers are not
    - Drivers can drive different cars
- In code:
  - Car would have Driver and Passenger as class fields
  - ...exactly like composition!

# Class Relationships (3)



- Association and Dependency
  - Indicates interaction between classes
    - Association = solid line, Dependency = dashed line
    - Difference is murky: UML is a *guide*, not a *law*
  - Used to show that one class invokes methods on another
    - ... but that there is no other relationship beyond this
    - With arrow, implies *unidirectional* (Car calls Weather, not vice-versa)
    - No arrow implies *bidirectional* (Car and Road call each other)
- In code: **Any way that a method call can be set up and made**
  - e.g., Weather object is passed as a parameter to a Car method
    - e.g., `Car.setAggressiveness(Weather currentConditions)`
  - e.g., Road has a class field of all Cars on that Road (aggregation?)



# Class Relationships (4)

- Inheritance

- “**is-a**” relationship

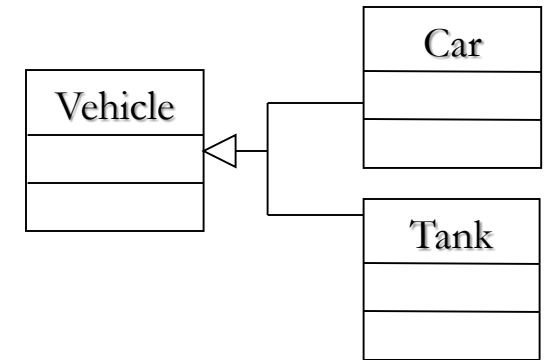
- Indicates one class is a sub-type of another class
    - Shown with an open triangle arrowhead beside super-type

- Implies the specialisation of the super-type

- Super-type synonyms: ‘parent’, ‘base’
  - Sub-type synonyms: ‘child’, ‘derived’

- In code: **During class declaration; syntax is language-specific**

- Python: `class Car(Vehicle):`
  - Java: `public class Car extends Vehicle`
  - C++/C#: `public class Car : Vehicle`



# Inheritance

- Inheritance is the ability of a new class or object to take on all of the properties of an existing class
  - i.e. the state and the functionality
- **Super Class:** The original class
- **Sub Class:** The new class which inherits all of the functionality of the super class
- The sub class can then:
  - Introduce additional state (class fields)
  - Modify the inherited functionality.
  - Introduce new functionality
    - i.e. more specialised
- The super class generally has less functionality than the sub class
  - i.e. more generalised

# Aggregation v's Inheritance

- An aggregation relationship is implied by the class field declarations
- An inheritance relationship is explicitly stated (given in brackets on the class definition)
- Note that BOTH relationships encapsulate the functionality of one class within another:
  - Any inheritance relationship can be re-expressed as an aggregation relationship and vice versa.
  - The choice is based upon which relationship is most appropriate.

# Class Responsibility

- Each class has a designated role or responsibility in the software system
- It may be that some classes have duplicated functionality
- This duplicated functionality can be removed and placed into a super class which the original classes inherit from
- It is important to ensure that a sub class never assumes the role of its super class
- If the sub class requires some super class functionality then it should call the appropriate super class method

# Super Class - Sub Class Communication

- Communication is one way:
  - Sub class calls super class methods but not the other way around
- The word *super* is used to refer to the super class
- `super()` by itself is a call to the super class' `__init__` method
- `super().methodName()` is a call to a public method in the super class
- Example:
  - In a super class there is a `toString()` method
    - `outStr = super().toString()`
  - The sub class `toString` method wishes to generate a string containing its own state plus the super class state:
    - `outStr = super().toString() + self.state`

# The Base Class

- All classes except one inherit from another class
- A special class, known as the *base class*, is *the only class that does not*
- In Python this base class is called *object*
- If no inheritance relationship is specified then it automatically inherits from the base class
  - Note: In Python 2, a class definition needed to state it inherited from object –  
`def class person(object)`

# Super Class / Sub Class

## Object Construction

- In order to construct a sub class object, a super class object must also be created
- The order of object construction is from the base class through to the sub class

# animals.py - Dog Class (Lecture 9)

```
class Dog():  
  
    myclass = "Dog"  
  
    def __init__(self, name, dob, colour, breed):  
        self.name = name  
        self.dob = dob  
        self.colour = colour  
        self.breed = breed  
  
    def printit(self):  
        print('Name: ', self.name)  
        print('DOB: ', self.dob)  
        print('Colour: ', self.colour)  
        print('Breed: ', self.breed)  
        print('Class: ', self.myclass)
```



# animals.py - Cat Class (Lecture 9)

```
class Cat():

    myclass = "Cat"

    def __init__(self, name, dob, colour, breed):
        self.name = name
        self.dob = dob
        self.colour = colour
        self.breed = breed

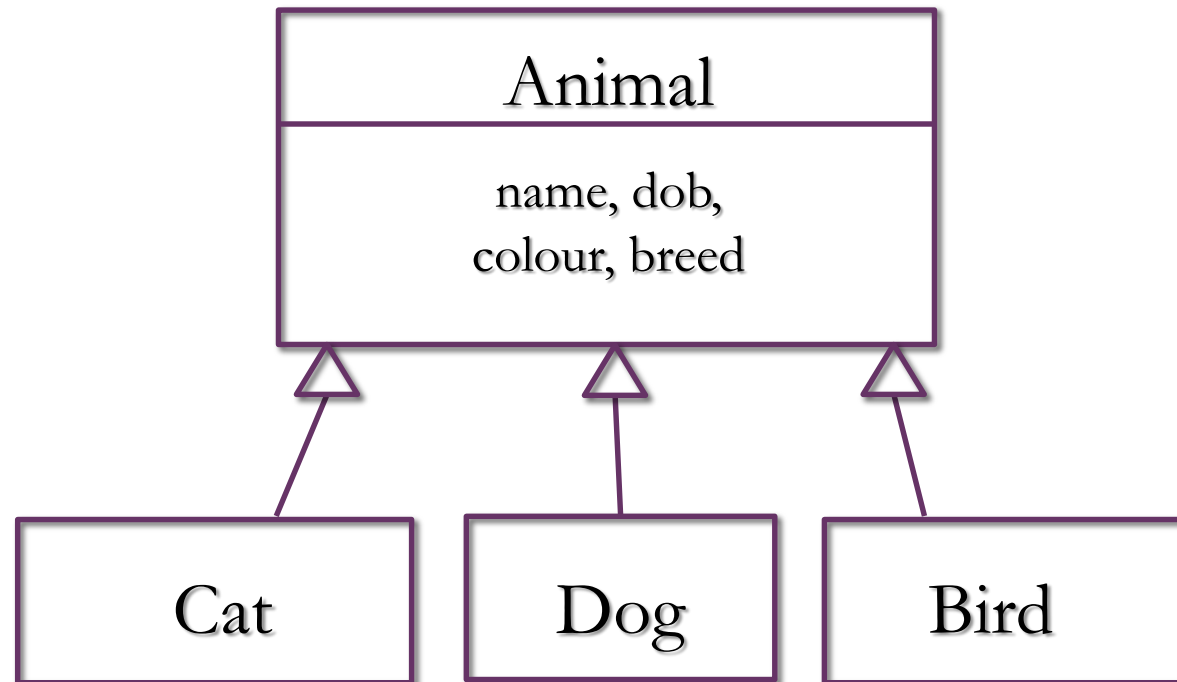
    def printit(self):
        print('Name: ', self.name)
        print('DOB: ', self.dob)
        print('Colour: ', self.colour)
        print('Breed: ', self.breed)
        print('Class: ', self.myclass)
```

# animals.py - Bird Class (Lecture 9)

```
class Bird():  
  
    myclass = "Bird"  
  
    def __init__(self, name, dob, colour, breed):  
        self.name = name  
        self.dob = dob  
        self.colour = colour  
        self.breed = breed  
  
    def printit(self):  
        print('Name: ', self.name)  
        print('DOB: ', self.dob)  
        print('Colour: ', self.colour)  
        print('Breed: ', self.breed)  
        print('Class: ', self.myclass)
```

# Example: Inheritance

- Repetition should be avoided if possible
- Cat, Dog and Bird are nearly identical
- Factor out the duplicated fields and methods...



# Example: animals.py

```
class Animal():

    myclass = "Animal"

    def __init__(self, name, dob, colour, breed):
        self.name = name
        self.dob = dob
        self.colour = colour
        self.breed = breed

    def __str__(self):
        return(self.name + '|' + self.dob + '|' + self.colour+'|'+self.breed)

    def printit(self):
        spacing = 5 - len(self.myclass)
        print(self.myclass.upper(), spacing*' ' + ': ', self.name, '\tDOB: ',
              self.dob, '\tColour: ', self.colour, '\tBreed: ', self.breed)
```

# Example: animals.py – magic!

```
class Dog(Animal):
```


```
    myclass = "Dog"
```

```
class Cat(Animal):
```

```
    myclass = "Cat"
```

```
class Bird(Animal):
```

```
    myclass = "Bird"
```



Just the differences  
between the **Animal**  
superclass and the  
subclasses

These changes would  
have no impact on  
Shelter.py or pets.py

# Polymorphism and Method Overriding

- An important aspect of inheritance for polymorphism is the ability to **override** methods of the base class
  - Consider passing a Tank to a method `void drive(Vehicle veh)`
  - A call to `veh.accelerate()` will actually call Tank's `accelerate()`
    - Which will behave differently to Car's `accelerate()`
- What is happening here?
  - Tank somehow becomes Vehicle. How?
  - What if you wanted to get back to Tank from Vehicle?
    - Since it really is a Tank, surely you can do it

# Overloading vs Overriding

- Overloading is when many methods share the same name but differ in their parameters
  - Constructors are a good example: default, alternate and copy constructor all have the same name, different parameters
    - Uniqueness is defined by name + parameter types
      - This is called the method's *signature*, or *prototype*
    - *e.g.*, Car(String model) and Car(int numSeats) are different
    - But: Car(String model) and Car(String ownerName) cannot be disambiguated - will cause compiler error
    - Note that return type is *not* part of the method signature
  - Most modern languages support overloading
    - C and Fortran are a couple that don't support overloading

# Overloading vs Overriding

- Overriding is where a method has exactly the same signature as a method in a *super/parent/base* class
  - *i.e.*, the child class is overriding the behaviour of the parent
  - Only applies to object-oriented languages, and all O-O languages support it
    - Overriding = specialisation, one of the cornerstones of O-O
- A method can be an overload *and* an override
  - Overloads the name of another method *in the **current** class*
  - Overrides the signature of a method *in the **parent** class*



# this, super keywords

- Keyword 'this' is a reference to the *current object*

e.g., 

```
public Tank clone() {  
    return new Tank(this); // Use copy constructor to make copy of ourselves  
}
```

- Keyword 'super' is a 'reference' to the current object's *parent class*
  - Use it to force a call to the parent class's code

e.g., 

```
public Tank(Tank otherTank) {    // Copy constructor  
    super(otherTank); // Call parent's copy constructor code first  
    // Now do our own copy constructor code... }
```

e.g., 

```
public void doSomething() {    // A method  
    super.doSomething();        // Call parent's doSomething() code  
    // Now do our own code... }
```

- **super** and **this** are relative to the current object/class
  - **this** = current object
  - **super** = current class's direct parent class

# Casting Between Types

- Changing from one type to another is called casting
  - You can also cast between numeric primitive types
    - *e.g.*, ints to floats and vice-versa, but not int to String.
    - C/C++ let's you cast *anything* - it's your problem if its wrong!

```
float fNum = 1.01;  
int iNum = (int) fNum;
```

← Cast by placing target data type in brackets

- Java (and pretty much every language) will implicitly do casts for you when it knows that the cast is 'safe'
  - Since Tank is-a Vehicle, casting Tank to Vehicle is safe

```
Tank t = new Tank();  
Vehicle v1 = t;  
Vehicle v2 = (Vehicle)t;
```

← Implicit cast

← Explicit cast, same result as implicit cast

- There's no need to explicitly do the casting here

# Casting Between Types (2)

- So when do you have to cast? And why?
  - When you are casting between numeric types
    - because loss of information can occur, *e.g.*, float 1.01 → int 1
  - When you are attempting to downcast to a derived class
    - *e.g.*, casting Vehicle to Tank is not safe since the compiler cannot be sure that the object (of known type Vehicle) is a Tank or not
      - Tank is-a Vehicle **does not mean** Vehicle is-a Tank!
- If you know the cast is OK you can do it explicitly
  - *e.g.*, You know that the Vehicle really is a Tank
  - Compiler then leaves it to run-time to try the cast
    - Fails at run-time with a ClassCastException if it's not a Tank

# Casting Between Types (3)

```
Vehicle v = new Tank();  
Tank t1 = v;  
Tank t2 = (Tank)v;  
Car c1 = (Car)v;
```

← Implicit cast is happening here  
← **Compiler error**  
← OK, and will work at run-time too  
← Will compile, but fails at run-time

- Some notes on casting
  - Primitives:
    - Casting from floats to ints will truncate the decimal places
    - Casting from ints to floats may lose some numerical precision
  - Classes
    - Object is a handy class to use for making general-purpose containers - simply contain an Object and you can contain *anything*
      - You have to explicitly cast back to the right class later though

# Checking Class Type

- Downcasting sounds a bit risky
  - What if you aren't totally sure of the object's true class?
    - Downcasting could cause a `ClassCastException`
    - Could catch this exception and try again, but that's ugly
  - Java provides you with a solution: `instanceof` keyword
    - Let's you check if object A is really an instance of class X

```
Vehicle v = new Tank();  
if (v instanceof Tank) {           ← Check if it really is a Tank  
    Tank t1 = (Tank)v;  
}
```

- **Warning:** try to limit your use of `instanceof` since it can be an indication of bad design and makes polymorphism redundant
  - Plus, if you are certain that the cast is OK, `instanceof` is a waste

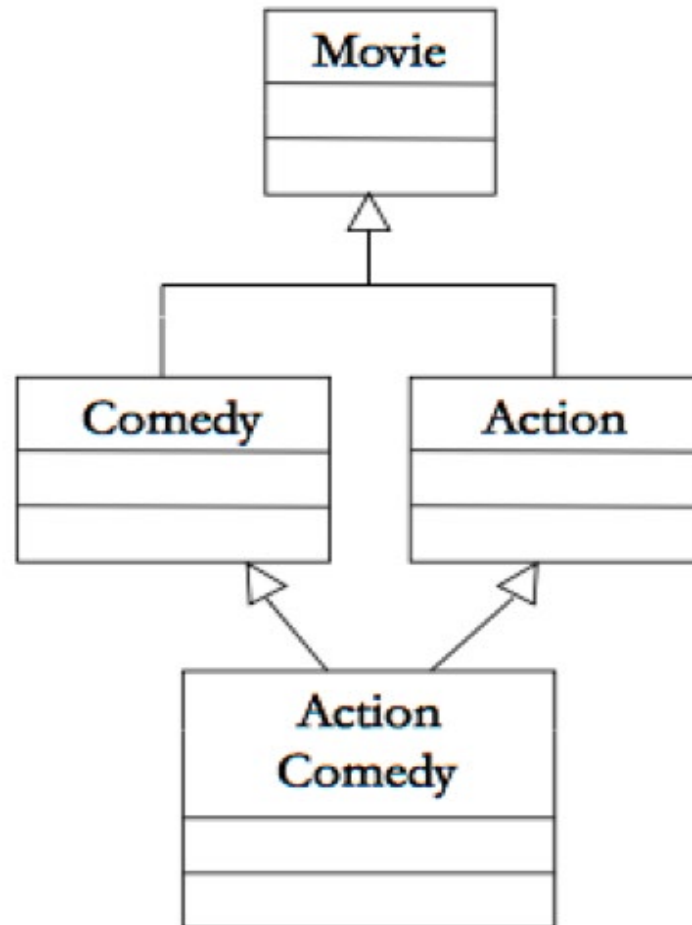
# Multiple Inheritance

- So what do we do if a class is required to inherit the state and functionality of more than one super class?
- So Tank “is-a” Vehicle
  - But Tank “is-an” Artillery as well, not just a Vehicle and Artillery is not always a Vehicle, so can’t put Artillery in between Tank and Vehicle
  - ie: Tank really has more than one base class
- One solution: allow multiple inheritance (eg: Python, C++)
  - Tank inherits from *both* Vehicle and Artillery

# Multiple Inheritance – Problems

- Theoretically, multiple inheritance is fine
- But in practice (in the code), things can get messy
- Say both Vehicle and Artillery define a method getSize()
  - If Tank does not override getSize(), which getSize() version should the compiler call? Vehicle's? Artillery's?
  - Worse, what if Artillery.getSize() refers to the size of the *shells* it fires, but Vehicle.getSize() refers to the *vehicle's* size?
- In more complicated inheritance hierarchies, you can even inherit from the same class more than once!
  - The next slide shows an example of this

# Multiple Inheritance - Example





# Interfaces (Java)

- Interfaces are used as a solution to resolve (some of) the problems with multiple inheritance
  - An interface is essentially an abstract class where:
    - All methods are abstract (ie: have no implementation)
    - All methods are public
    - No class fields exist
- In other words, an interface class only defines a set of public methods that its child classes must implement
  - Note that interfaces cannot have a constructor
    - There's nothing to construct, so what would be the point?
  - Interfaces can inherit from (extend) other interfaces, but do not have to (unlike classes, which extend at least Object)

# Interfaces and Multiple Inheritance (Java)

- Many multiple inheritance issues can then be resolved
  - Allow inheritance from as **many** interfaces as required
    - Interface inheritance
  - BUT only allow inheritance from a **single** class, which includes abstract classes
    - Implementation inheritance
- Why does this help?
  - Because interfaces cannot have any code
  - Thus there is never any confusion as to which base class's method should be invoked - there is only ever one base class with an implementation (all others are interfaces)

# Interfaces and Multiple Inheritance (Java)

- Interfaces are not a magic cure-all
  - *e.g.*, If Vehicle and Artillery are both made into interfaces, but `getSize()` has different meanings for both:
    - Tank still can't properly choose how to override `getSize()`
    - C# has the ability to define different methods, one per interface
  - *e.g.*, Action Comedy
    - Action and Comedy aren't abstract, and so can't be interfaces
    - Could make *all* movie genres into interfaces, and have separate implementation classes inheriting from these. Messy!
  - Limits code reuse potential
    - Interfaces have no implementation (code) to reuse!

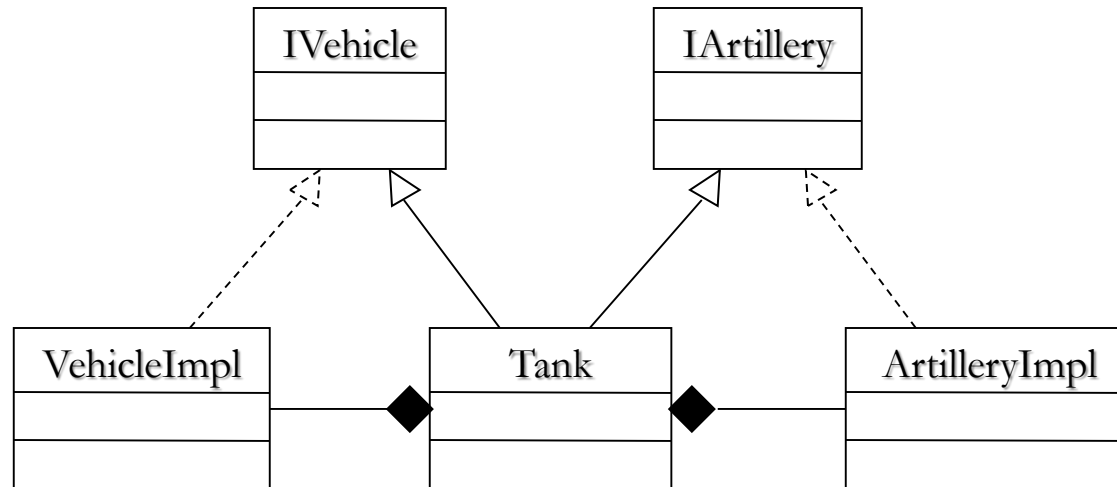
# Interfaces and Multiple Inheritance (Java)

- Interfaces are not a magic cure-all
  - *e.g.*, If Vehicle and Artillery are both made into interfaces, but `getSize()` has different meanings for both:
    - Tank still can't properly choose how to override `getSize()`
    - C# has the ability to define different methods, one per interface
  - *e.g.*, Action Comedy
    - Action and Comedy aren't abstract, and so can't be interfaces
    - Could make *all* movie genres into interfaces, and have separate implementation classes inheriting from these. Messy!
  - Limits code reuse potential
    - Interfaces have no implementation (code) to reuse!

# Emulating Multiple Impl Inheritance

- Ideally Tank would inherit from Vehicle and Artillery
  - ...and both would have code that Tank can reuse
    - *i.e.*, they are not interfaces, probably abstract classes instead
- The aforementioned issues with M.I. are in our way
  - But we can *emulate* M.I. with interfaces and composition
    - Have Tank inherit from interface **IVehicle**
    - Have Tank compose with (contain) a class VehicleImpl that implements all the would-be-non-abstract methods of **IVehicle**
      - VehicleImpl might also inherit from **IVehicle**, but will have to bomb out on any truly-abstract method - a bit messy
    - Have Tank 'delegate' calls to equivalent methods in VehicleImpl
      - VehicleImpl code can then be shared (re-used) with other classes
    - Then do the same with **IArtillery**

# Emulating M.I. - Example



# Class vs Abstr Class vs Interface Inheritance

- Inherit from classes...
  - ...when you need to specialise behaviour of existing class
- Inherit from abstract classes...
  - ...where a lot of the code in derived classes is common among most/all of the derived classes
  - The abstract class is then a ‘repository’ for shared code
- Use interfaces and composition+delegation...
  - ...everywhere else
    - It avoids wasting your precious single base class
    - It also helps you get around integrating with or reusing existing classes - inherit from one, compose+delegate with others

# Interfaces in Code

- Naming:
  - A prefixed capital 'I' is common for interfaces, eg: IVehicle
- In code:
  - Declaring: Almost identical to declaring a class
    - Java: `public interface IVehicle { ... methods here ... }`
    - C#: `public interface IVehicle { ... methods here ... }`
  - Inheriting from:
    - Java: `public class Tank implements IVehicle, IArtillery`
    - C#: `public class Car : IVehicle, IArtillery`
  - Can use extends and implements keywords together:
    - `public class Tank extends MilitaryObject implements IVehicle, IArtillery`



# ABSTRACT DATA TYPES

---

# SortedList – an Abstract Data Type

- Last week we saw the value of sorting, and some sorting algorithms
- Our searching will be faster if the data is sorted
- How do we maintain a sorted list if the data is changing?
- Over time we may need to insert and delete values
- Create a class SortedList holding the list and the current number of elements
- The main operations are:

**find – insert - delete**

# SortedList - Class Diagram (UML)

## **SortedList**

- theSortedList : array of integers
- numElements : integer

- + find(int key): int
- + insert(int key): none
- + delete(int key): none

# find

- assume theSortedList and numElements are classfields

Submodule: find (AKA linear search)

Import: key (item to find)

Export: location (index)

Assertion: returns the location of key if it exists in the array,  
otherwise throws an exception

```
location=0, found = false
```

```
DO
```

```
    IF theSortedList[location].equals <-- key
```

```
        found = true
```

```
    ELSE
```

```
        increment location
```

```
WHILE NOT found AND location < numElements
```

```
IF NOT found
```

```
    throw appropriate exception
```

# insert

- three scenarios

1. **End of list of values**

- position = element [numElements]
- easy!

```
IF theSortedList is not full           ← throw exception if it is!
    theSortedList [numElements] = insertValue
    increment numElements
ENDIF
```

1. **Beginning of list of values**

- element [0]

```
IF theSortedList is not full           ← throw exception if it is!
    FOR ii= numElements, ii>0, decrement ii  ← Shuffle elements away to make room
        theSortedList[ii] = theSortedList[ii-1]
    ENDFOR
    theSortedList[0] = insertValue
    increment numElements
ENDIF
```

Submodule: insert

Import: key (item to insert)

Export: None

Assertion: inserts key/value into array at correct, sorted position, otherwise throws an exception

# insert

## 3. Somewhere inside list of values

- Need to search for position, then insert
- Array needs to shuffle down to make space

```
position=0
IF theSortedList is not full                                ← throw exception if it is!
    WHILE insertValue < theSortedList[position] AND position < numElements
        increment position
    ENDWHILE

    FOR ii=numElements, ii>position, decrement ii           ← Shuffle elements away
        theSortedList[ii] = theSortedList[ii-1]
    ENDFOR
    theSortedList[position]=insertValue
    increment numElements
ENDIF
```

# delete (remove)

- need to ensure array is not empty!

- throw exception if it is

- three scenarios

1. End of list of values

- Element [n-1] is deleted
- Decrement count.

2. Beginning of list of values

- Element [0] is deleted
- Starting from element [1], shuffle the rest of the elements down by one, overwriting element [0].
- Decrement count.

3. Somewhere else in list

- Element[x?]
- Find the element to delete.
- Starting from the next element, shuffle the rest of the elements down by one, overwriting the element to delete.
- Decrement count.

Submodule: delete

Import: key (item to insert)

Export: None

Assertion: deletes key/value from array, otherwise throws an exception

# Time Complexity – SortedList

| Operation | Best Case | Average Case | Worst Case |
|-----------|-----------|--------------|------------|
| find      | $O(1)$    | $O(N)$       | $O(N)$     |
| insert    | $O(N)$    | $O(N)$       | $O(N)$     |
| delete    | $O(N)$    | $O(N)$       | $O(N)$     |

Each operation needs to do a "find", then a possible shuffle  $O(N)$



# Data Structures

- Arrays are a type of *data structure*
  - They define how to organise data in memory
  - In particular, arrays store a set of elements in a single contiguous block of memory, accessed via an index
- Data structures such as arrays can be useful as they are, but they aren't always a perfect fit
  - Many applications need to access data differently to the array's 'index-update' approach
    - e.g., an order processing queue: take from front, add to rear
  - Problem: an array is really *how a computer operates*
    - RAM is just one long 1D array (same with disk storage)

# Abstract Data Types

- So there can be a gap between the data structure (how it works) and the *usage* of that structure
- Abstract Data Types are there to define behaviour
  - ADT: a set of methods that provide access to data in a way that is natural for the application
  - How the methods manipulate the underlying data structure to achieve this is not the app's problem
    - Even the data structure used is hidden!
  - ADTs make developing applications much easier
    - Write the ugly details once and wrap it all in nice methods
    - Lets you later concentrate on the application logic rather than the details of manipulating the data structure

# Abstract Data Types as Objects

- ADTs are defined in terms of operations
- Objects bundle state and operations together
- Our objects (classes) must include
  - Code to implement all ADT operations
  - Instance variables to support the required state (e.g. array of data, count)
  - Methods for initialising the objects
  - Support methods, e.g. display( )
  - Validation and exception handling throughout
- We may choose different internal implementations:
  - Data types and structures (e.g. arrays, lists, trees)
  - Algorithms (e.g. sorting, searching, traversing)

# STACKS AND QUEUES

---

# Stacks and Queues

- Two very common ADTs are stacks and queues
  - Queue: elements taken out in the order they were added
    - FIFO: first-in, first-out (although not all queues are FIFO queues)
  - Stack: data elements are taken out in *reverse* order
    - LIFO: last-in, first-out
  - Elements *must* be taken out in the appropriate order: you can't jump in and grab the 5<sup>th</sup> element
- Such processing occurs a lot in the real world
  - And we often need to model such processes in software
- **But:** arrays aren't necessarily best for implementing these ADTs

# Queue vs Array

- Consider the behaviour of a queue vs an array:
  - Nothing stops you from accessing array element [5]
    - But a queue should only take the first element each time
  - If you take the first array element [0], element [1] doesn't automatically move to position [0]
    - So then you have to remember that the 'new-first' element is [1],
    - or shuffle all the elements up by one yourself
- Solution: use methods to make the array behave like a queue
  - Just because it's messy doesn't mean it's impossible
    - ...but it means we only have to **CODE AND TEST IT ONCE!**
  - If we code it right, using it in the application will simplify (and clarify) the rest of the code enormously

# Stacks

- Let's start with stacks, because they are easier!
- A stack is an ADT that implements a LIFO list
  - Think of a stack of plates – add to top, take from top
  - Some example applications for stacks:
    - Converting a character **string** into an **int** (e.g., “10” → 10)
    - Storing information for method calls
    - Evaluating a mathematical expression ( We'll see later on)
- Since it's an ADT, we'll first talk about *what* a stack's behaviour is
  - Then we will discuss *how* to implement a stack
    - In particular: with an array data structure (this time)

# Stack Methods

- Being LIFO, a stack has a few obvious methods, with standard names that everyone recognises:
  - `push()` – add a new item to the top of the stack
  - `pop()` – take the top-most item from the stack
  - `top()` – look at the top-most item, but leave it on the stack
    - Synonym: `peek()`
  - `isEmpty()` – check if the stack is empty
- There are also extra methods that often appear
  - `isFull()` – checks if the stack is full
    - Arrays can get full, but some data structures don't have this issue
  - `count()` – number of elements in the stack
    - Synonyms: `size()`, `numElements()` (not as standardised!)



# Stack Implemented with an Array

- Java and Python have built-in classes for stacks, but we'll develop our own DSAShStack to illustrate the concept
  - DO NOT USE BUILT-IN DATA TYPES AND ALGORITHMS IN DSA
  - Let's create a stack of double values to hold numbers
- The only data structure we know (so far) for storing sets of data is the array ... so we'll use arrays
- How are we going to do it?
  - Look for similarities that we can exploit
  - Consider: A stack grows and shrinks on *one side*
  - Similarly, array elements start at [0], and can be added to / removed from the end until the array capacity is reached

# Stacks with Arrays

- So, if we make the *top* of the stack be the *back* of the array, we can grow/shrink without much hassle
  - Counter-intuitive, but simplifies the code a lot!
- The idea is to keep track of the count of elements in the array
  - The element at `[count - 1]` is then the top of the stack
    - `- 1` because arrays are zero-based in Java/Python, remember!
  - New items then get stored in slot `[count]`
    - `[count-1]` is the top, so `[count]` is the next unused slot
  - When `count == array.length`, the stack is Full

# Stack - Pseudocode

```
Class DSASStack
Class fields : stack (double array), count (integer)
Class constant : DEFAULT_CAPACITY ← 100
```

```
Default constructor
  alloc stack array with DEFAULT_CAPACITY elements
  count ← 0
```

```
Alternate constructor IMPORT maxCapacity (integer)
  alloc stack array with maxCapacity elements
  count ← 0
```

```
ACCESSOR getCount IMPORT none EXPORT count
```

```
ACCESSOR isEmpty IMPORT none EXPORT empty (boolean)
  empty ← (count = 0)
```

```
ACCESSOR isFull IMPORT none EXPORT full (boolean)
  full ← (count = stack length)
```

<continued next slide>

# Stack - Pseudocode (cont.)

```
MUTATOR push IMPORT value EXPORT none
  IF isFull() THEN
    ABORT                                ← ie: throw an exception
  ELSE
    stack[count] ← value
    count ← count + 1
  ENDIF
```

```
MUTATOR pop IMPORT none EXPORT topVal
  topVal ← top()
  count ← count - 1
```

```
ACCESSOR top IMPORT none EXPORT topVal
  IF isEmpty() THEN
    ABORT
  ELSE
    topVal ← stack[count - 1]
  ENDIF
```

# Application: Palindrome

- How can we check if a string (or number) is a palindrome?
- Need to check if it's the same forward and backward.
- We can achieve this with a stack...

```
IMPORT: inString
EXPORT: match
create a new palStack
FOR ch ← 0 TO inString.length -1 DO
    palStack.push ← ch
ENDFOR

pos = 0
match = TRUE

WHILE match AND NOT palStack.isEmpty
    match = inString[pos] == palStack.pop
    pos = pos + 1
ENDWHILE
```

# Application: ReadInt

- In the lecture on recursion we saw that the system stack can be used to convert characters read from the keyboard to an integer.
- We can also achieve this with our own stack.

```
create a new intStack
ch = readChar
WHILE '0' <= ch <= '9'
    digit = ch - '0'
    intStack.push<-- digit
    ch = readChar
ENDWHILE

value = 0
powerOfTen = 1

WHILE NOT intStack.isEmpty
    digit = intStack.pop
    value = value + digit * powerOfTen
    powerOfTen *= 10
ENDWHILE
```

# Application: Evaluation of Maths Equations

- Stacks *really* become useful for non-obvious tasks
  - Evaluation of maths expressions is one of those tasks
- The problem:
  - We normally see equations in the form:
$$(10.3 * (14 + 3.2)) / (5 - 2 * 3)$$
  - There are many precedence rules that need to be followed
    - BIMDAS or BOMDAS
    - Makes it hard to write code to solve it in the right order

# Infix to Postfix

- Solution: Re-order the equation so that higher precedence operations come before lower ones
  - Plus we get rid of brackets, even nested brackets
  - Then we just need to read it from left-to-right
- How?
  - Normal equations are in what is called 'infix' notation
    - Unfortunately it's not possible to rewrite equations in infix to get rid of precedence ordering and brackets. Consider:  
Normal:  $(10.3 * (14 + 3.2)) / (5 + 2 - 4 * 3)$   
Left-to-Right:  $14 + 3.2 * 10.3 / -4 * 3 + 5 + 2$  (ie: no BIMDAS)
  - Close, but the  $10.3 / -4$  is wrong – we needed to 'postpone' evaluating it until after the  $+ 2$ . But with infix we can't postpone



# Postfix

- Solution: use a different notation, **postfix**
  - Put the operator *after* the operands it applies to (the 'post')
  - Each operator then applies to the two operands that precede the operator
- How does this help?
  - You only evaluate operands once you see an operator
    - Before that, you just keep adding operands to a pile
    - Since the operator must be applied to the *last* two operands (LIFO), your 'pile' is in fact a **stack**

# Infix vs Postfix Examples

- The original equation in Postfix:

Infix:  $(10.3 * (14 + 3.2)) / (5 + 2 - 4 * 3)$

Postfix:  $10.3\ 14\ 3.2\ +\ *\ 5\ 2\ +\ 4\ 3\ *\ -\ /\$

- Some simpler examples:

| Infix                         | Postfix                           |
|-------------------------------|-----------------------------------|
| $3 * 4$                       | $3\ 4\ *$                         |
| $2 - 4 + 3$                   | $2\ 4\ -\ 3\ +$                   |
| $4 + 2 * 3$                   | $4\ 2\ 3\ *\ +$                   |
| $(4 + 2) * 3$                 | $4\ 2\ +\ 3\ *$                   |
| $((2 - 3) / 4 * (1 + 9)) * 2$ | $2\ 3\ -\ 4\ /\ 1\ 9\ +\ *\ 2\ *$ |

# Postfix Properties

- Points to note:
  - The order of the operands is left **unchanged**
  - Operators are listed in **precedence order**
    - ... even the effect of brackets has been taken into account
  - Equal-precedence operators are kept in the infix order
    - left to right associativity
      - e.g.,  $2 - 4 + 3 \rightarrow 2\ 4 - 3 +$  NOT  $2\ 4\ 3 + -$
      - Reason:  $2 - 4$  is in fact  $2 + (-4)$ , so we *must* keep the -ve sign related to the 4:  $2 - 4 \neq 4 - 2$
      - $2\ 4\ 3 + -$  is actually postfix for  $2 - (4 + 3)$
    - Same reasoning applies to  $\backslash$ :  $A \backslash B \neq B \backslash A$
    - $+$  and  $*$  aren't so problematic, since  $A + B = B + A$

# Evaluating Postfix

- Evaluating postfix expressions will give some more insight into why it all works
  - We'll discuss infix  $\rightarrow$  postfix conversion a little later
    - ... because it's harder!
- Unsurprisingly, we use a stack in the evaluation
  - Push operands onto stack until an operator is encountered
  - Pop off last two operands and apply the operator to them
    - Apply the operator *in-order*, not LIFO order (important for  $-$ ,  $/$ )
  - Push the result back on the stack ready for the next op
  - When no more operands/operators are left in the postfix, the answer is the (single) value remaining on the stack

# Postfix Evaluation Example

Infix:  $(10.3 * (14 + 3.2)) / (5 + 2 - 4 * 3)$

Postfix: 10.3 14 3.2 + \* 5 2 + 4 3 \* - /

| PFix  | Eval Stack Contents | What's Happening?                               |
|-------|---------------------|-------------------------------------------------|
| 10.3  | 10.3                | <push 10.3>                                     |
| 14    | 10.3 14             | <push 14>                                       |
| 3.2   | 10.3 14 3.2         | <push 3.2>                                      |
| +     | 10.3 17.2           | <2 pops> $\rightarrow 14 + 3.2$ , <push ans>    |
| *     | 177.16              | <2 pops> $\rightarrow 10.3 * 17.2$ , <push ans> |
| 5     | 177.16 5            | <push 5>                                        |
| 2     | 177.16 5 2          | <push 2>                                        |
| +     | 177.16 7            | <2 pops> $\rightarrow 5 + 2$ , <push ans>       |
| 4     | 177.16 7 4          | <push 4>                                        |
| 3     | 177.16 7 4 3        | <push 3>                                        |
| *     | 177.16 7 12         | <2 pops> $\rightarrow 4 * 3$ , <push ans>       |
| -     | 177.16 -5           | <2 pops> $\rightarrow 7 - 12$ , <push ans>      |
| /     | -35.432             | <2 pops> $\rightarrow 177.16 / -5$ , <push ans> |
| <end> | -35.432             | <pop> $\rightarrow$ Final answer                |

# Infix to Postfix Conversion

- Converting infix to postfix *also* uses a stack
  - Postfix needs to re-arrange operators into the right place
  - So we need to 'hold on' to operators until we reach the right point in the equation to insert them back in
    - Remember that operands don't change their order
  - The method behind this is to hold back an operator until we see an equal-or-lower-precedence operator
    - If the new operator is higher precedence, we have to put it 'on top' of the other operator (in a stack), since it takes precedence
  - Brackets are an extra wrinkle
    - Approach: treat sub-equations in brackets as if they were isolated from the rest of the equation (because they are!)

# Infix to Postfix Conversion: Algorithm

```
postfix ← empty
WHILE infix has more terms DO
    term ← ParseNextTerm()

    IF (term = '(') THEN
        opStack.push('(')

    ELSE IF (term = ')') THEN
        WHILE (opStack.top ≠ '(') DO
            postfix ← postfix + opStack.pop
        ENDWHILE
        opStack.pop

    ELSE IF (term = '+' ) OR (term = '-' ) OR (term = '*' ) OR (term = '/' ) THEN
        WHILE (NOT opStack.isEmpty) AND (opStack.top ≠ '(') AND
            (PrecedenceOf(opStack.top) ≥ PrecedenceOf(term)) DO
            postfix ← postfix + opStack.pop
        ENDWHILE
        opStack.push(term)

    ELSE
        postfix ← postfix + term
    ENDIF
ENDWHILE

WHILE (NOT opStack.isEmpty) DO
    postfix ← postfix + opStack.pop
ENDWHILE
```

NOTE: Methods in red must also be implemented, but are fairly straightforward tasks

- ← Extract next term (operator, operand) from infix eqn
- ← '(' gets put straight onto the stack
- ← Find corresponding '('
- ← Pop remaining operators for the bracketed sub-equation
- ← Pop the '(' and discard it
- ← Move higher/equal precedence ops to postfix eqn
- ← Always put the new operator onto the stack
- ← Term must be an operand if it isn't an operator
- ← Add operand to postfix equation
- ← Pop any remaining operators from the stack

# Infix to Postfix Example

Infix:  $(10.3 * (14 + 3.2)) / (5 + 2 - 4 * 3)$

Postfix:  $10.3\ 14\ 3.2\ +\ *\ 5\ 2\ +\ 4\ 3\ *\ -\ /\$

| Infix | Postfix So Far                   | Operator Stack |
|-------|----------------------------------|----------------|
| (     |                                  | (              |
| 10.3  | 10.3                             | (              |
| *     | 10.3                             | ( *            |
| (     | 10.3                             | ( * (          |
| 14    | 10.3 14                          | ( * (          |
| +     | 10.3 14                          | ( * ( +        |
| 3.2   | 10.3 14 3.2                      | ( * ( +        |
| )     | 10.3 14 3.2 +                    | ( *            |
| )     | 10.3 14 3.2 + *                  | <empty>        |
| /     | 10.3 14 3.2 + *                  | /              |
| (     | 10.3 14 3.2 + *                  | / (            |
| 5     | 10.3 14 3.2 + * 5                | / (            |
| +     | 10.3 14 3.2 + * 5 2              | / ( +          |
| 2     | 10.3 14 3.2 + * 5 2              | / ( +          |
| -     | 10.3 14 3.2 + * 5 2 +            | / ( -          |
| 4     | 10.3 14 3.2 + * 5 2 + 4          | / ( -          |
| *     | 10.3 14 3.2 + * 5 2 + 4          | / ( - *        |
| 3     | 10.3 14 3.2 + * 5 2 + 4 3        | / ( - *        |
| )     | 10.3 14 3.2 + * 5 2 + 4 3 * -    | /              |
| <end> | 10.3 14 3.2 + * 5 2 + 4 3 * - /\ | <empty>        |



# Postfix Conversion 'Checklist'

- Things to keep in mind:
  - Don't forget to write down the brackets in the infix!
  - New operators ALWAYS go onto the stack
    - They *never* get put directly onto the postfix expression
    - The only question is whether to first pop the operator that is *already on the stack* off to the postfix expression
  - Brackets NEVER appear in the postfix
    - And closing brackets never appear in the operator stack – they are only markers to indicate the end of the sub-equation
  - Remember to pop off any remaining operators at the end of each sub-equation or at the end of the full equation

# FIFO Queues

- A FIFO queue is an ADT implementing a FIFO list
  - Other kinds of queues aren't FIFO, eg: priority queue
- Examples of where FIFO queues are needed
  - Bank transactions: processed in the order they are made
  - Customer orders: first come, first served

# Queue Methods

- Queues (FIFO or otherwise) have the following methods
  - Note: naming isn't as standardised as it is with stacks
  - enqueue() – add item to the queue
    - FIFO queues add to the end, priority queues insert in priority order
    - Synonyms: add(), insert()
  - dequeue() – take item from the front of the queue
    - Synonyms: remove(), delete()
  - peek() – check the front item, but don't take it off
    - Synonyms: front()
  - isEmpty() – check if the queue is empty
  - isFull() – check if the queue is full. Optional
  - count() - number of elements in the queue. Optional

# FIFO Queue with an Array

- Unlike stacks, queues grow on one side (the end) and shrink on the other (the front)
  - No synergies with arrays to be taken advantage of here!
- Two options are available:
  - Shuffle queue elements forward when front is dequeued
    - Exactly like a real-world queue, like at the bank
  - Leave elements as-is and change which index is 'front'
    - *i.e.*, dequeued indexes are no longer used
    - Circular queue: allow the queue to cycle around the array, so that previously-dequeued indexes can be re-used

# 'Shuffling' vs Circular Queues

- Time Efficiency:
  - **Shuffling**: every dequeue must move N elements up by 1
  - **Circular**: Only need to adjust front index – much faster
- Space Efficiency:
  - Both have same space usage: circular queues can just start at idx [5], go through [length-1] and wrap around to end at [4].
  - But both still have a maximum size (due to fixed-size array)
- Code Complexity:
  - **Shuffling**: easy to understand, code, and maintain
  - **Circular**: Dealing with the wrap-around can be tricky – simplify it by storing the count as well as start/end indexes

# FIFO Queue – Pseudocode (Shuffling)

```
Class DSAQueue
Class field : queue (double array), count (integer)
Class constant : DEFAULT_CAPACITY ← 100
```

Default constructor

```
// implement this yourself
```

Alternate constructor **IMPORT** maxCapacity (integer)

```
// implement this yourself
```

**ACCESSOR** getCount **IMPORT** none **EXPORT** count

**ACCESSOR** isEmpty **IMPORT** none **EXPORT** empty (boolean)

```
// implement this yourself
```

**ACCESSOR** isFull **IMPORT** none **EXPORT** full (boolean)

```
// implement this yourself
```

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# FIFO Queue – Pseudocode (cont.)

```
MUTATOR enqueue IMPORT value EXPORT none  
  // implement this yourself
```

```
MUTATOR dequeue IMPORT none EXPORT frontVal  
  // implement this yourself
```

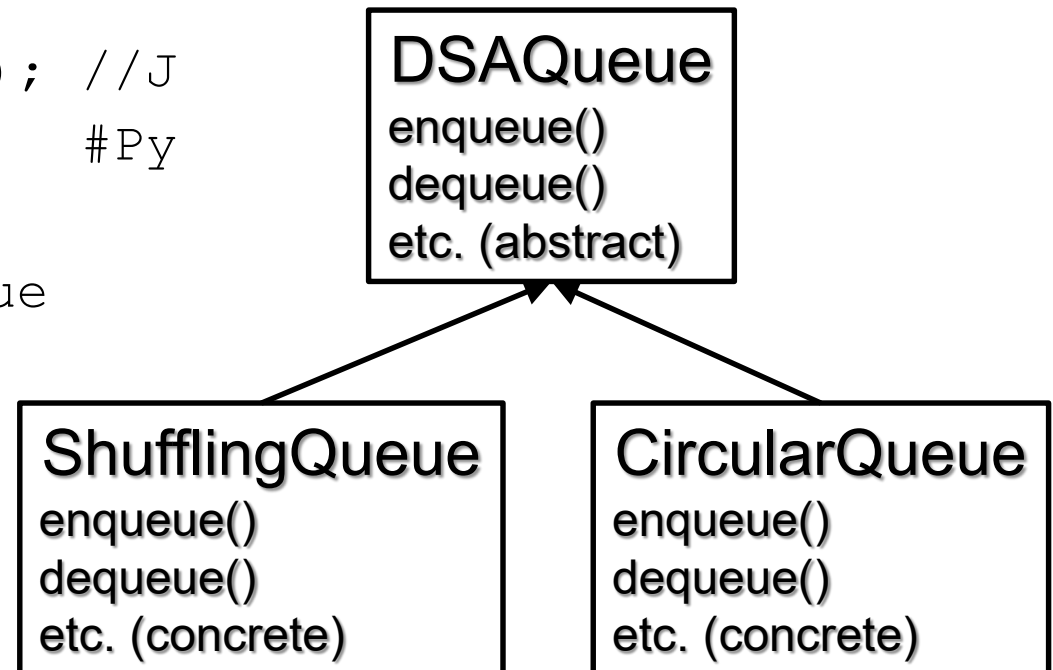
```
ACCESSOR peek IMPORT none EXPORT frontVal  
  // implement this yourself
```

# FIFO Queues - Polymorphism

- We can implement queues as shuffling or circular queues
- Using polymorphism, we can minimise changes required
  - Switch between implementations by changing one line of code

```
myQ = new ShufflingQueue( ); //J  
myQ = ShufflingQueue( )      #Py
```

```
// use methods from DSAQueue  
myQ.enqueue(200)  
myQ.peak( )
```





# Next Week

- Linked lists
- Iterators