



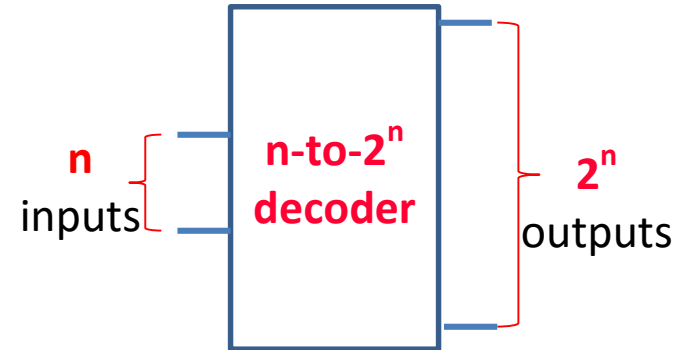
Digital Logic Design

Decoder & Encoder



Decoder

A **n-to- 2^n decoder** takes an **n-bit** input and produces **2^n** outputs. The **n** inputs represent a binary number that determines which of the **2^n** outputs is *uniquely* true.



Example:

- Reception counter: When you reach an Academic Institute
 - Receptionist asks: Which Dept. to go?
 - Based on your Specific answer, Receptionist redirects you to the specific building.

The job of the Decoder is to **Decode!**

-It knows what to do for a fixed question.

Use:

- Memory addressing
 - Address to a particular location.

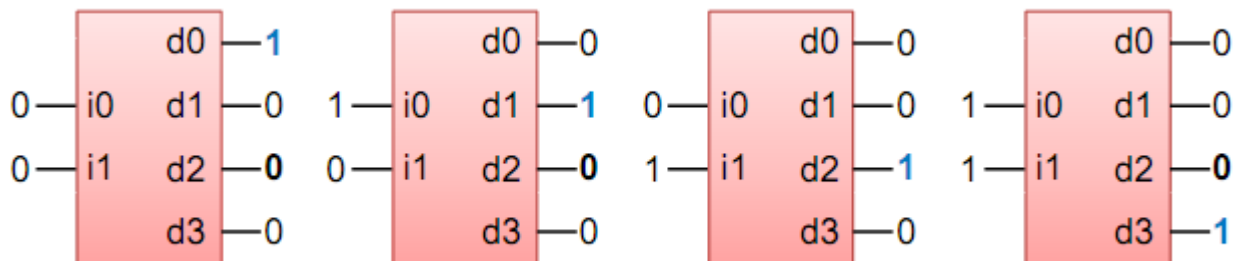
2-to-4 decoder

- A 2-to-4 decoder operates according to the following truth table.
 - The 2-bit input is called $S_1 S_0$, and the four outputs are $Q_0 Q_1 Q_2 Q_3$.
 - If the input is the binary number i , then output Q_i is uniquely true.

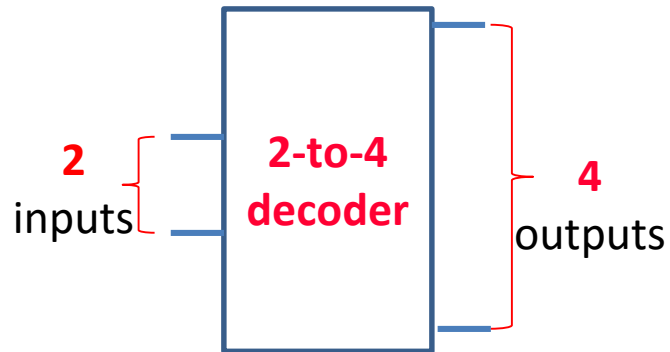
S_1	S_0	Q_0	Q_1	Q_2	Q_3
0	0	1	0	0	0
0	1	0	1	0	0
1	0	0	0	1	0
1	1	0	0	0	1

For instance, if the input $S_1 S_0 = 10$ (decimal 2), then output Q_2 is **true**, and Q_0, Q_1, Q_3 are all **false**.

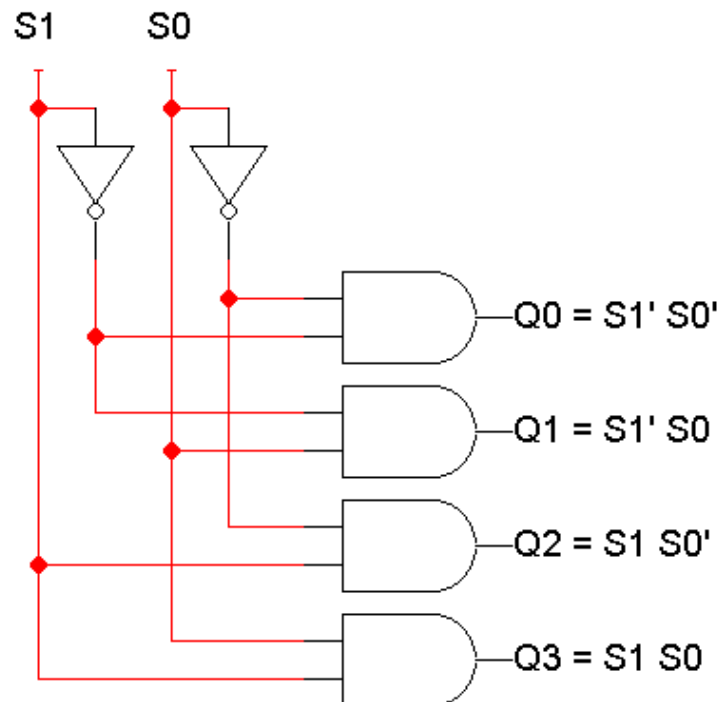
This circuit “decodes” a binary number into a “one-of-four” code.



Logic diagram of a 2-to-4 decoder



S1	S0	Q0	Q1	Q2	Q3
0	0	1	0	0	0
0	1	0	1	0	0
1	0	0	0	1	0
1	1	0	0	0	1



$$Q0 = S1' S0'$$

$$Q1 = S1' S0$$

$$Q2 = S1 S0'$$

$$Q3 = S1 S0$$

- Decoding ONLY a specific sequence:

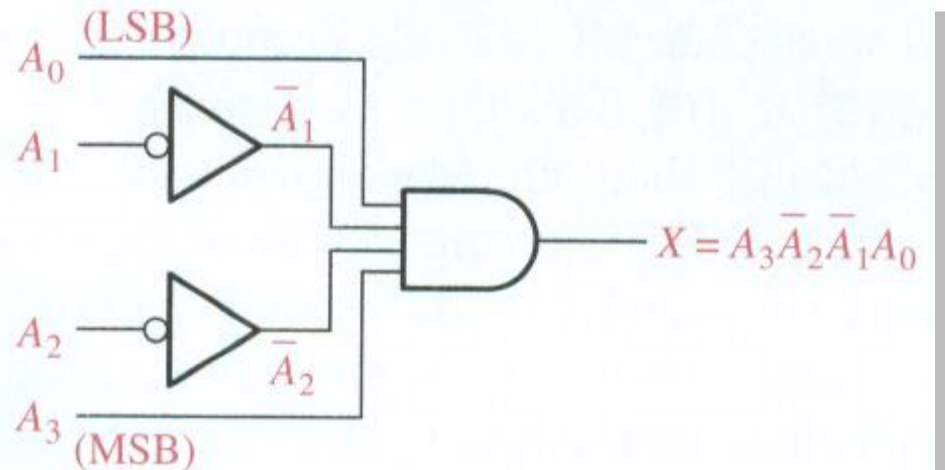
The output is 1 only when:

$$A_0 = 1$$

$$A_1 = 0$$

$$A_2 = 0$$

$$A_3 = 1$$



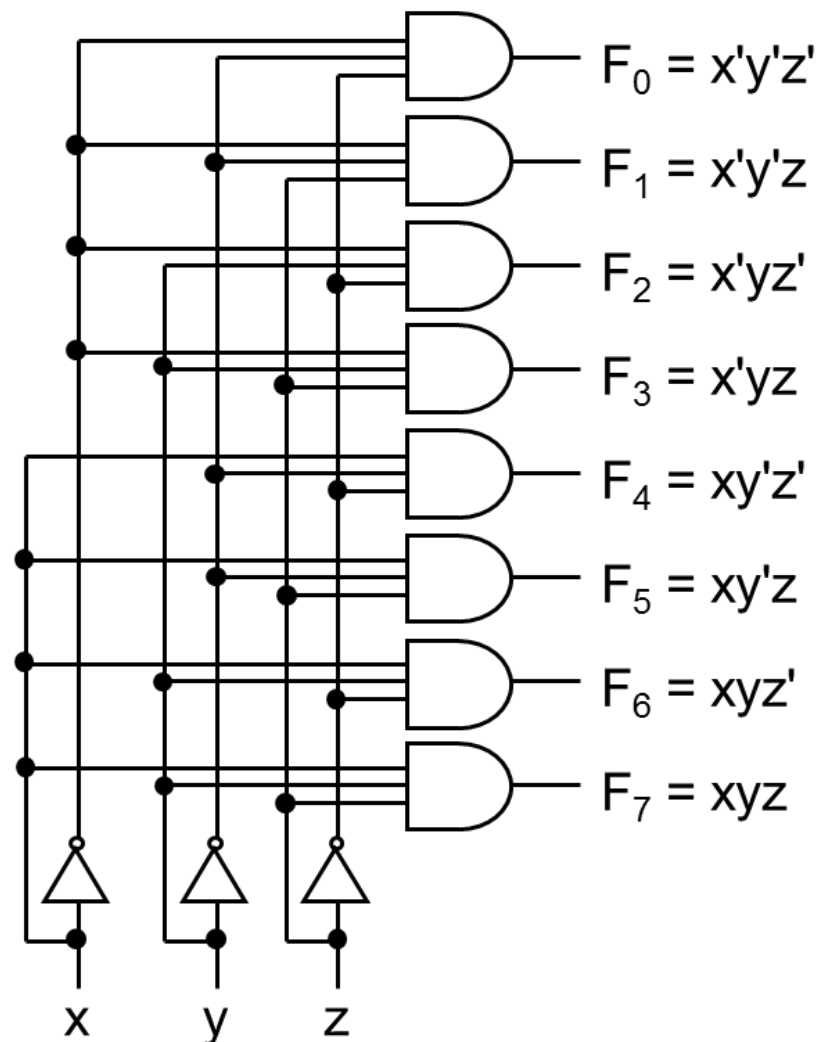
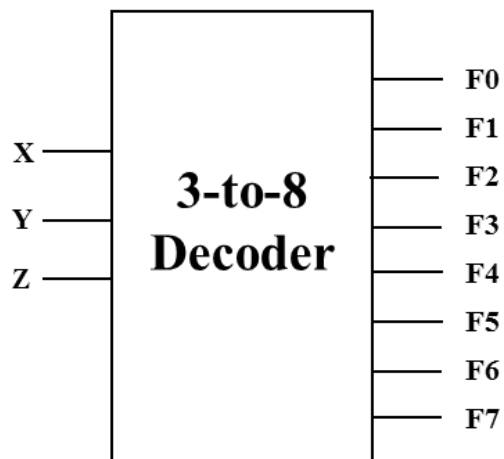
Use:

- 1) Encryption system,
- 2) Counter decoding...etc.

3-to-8 Binary Decoder

Truth Table:

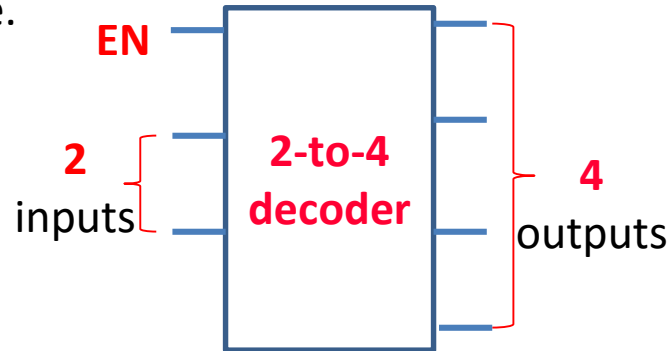
x	y	z	F ₀	F ₁	F ₂	F ₃	F ₄	F ₅	F ₆	F ₇
0	0	0	1	0	0	0	0	0	0	0
0	0	1	0	1	0	0	0	0	0	0
0	1	0	0	0	1	0	0	0	0	0
0	1	1	0	0	0	1	0	0	0	0
1	0	0	0	0	0	0	1	0	0	0
1	0	1	0	0	0	0	0	1	0	0
1	1	0	0	0	0	0	0	0	1	0
1	1	1	0	0	0	0	0	0	0	1



[illegible]

Enable inputs

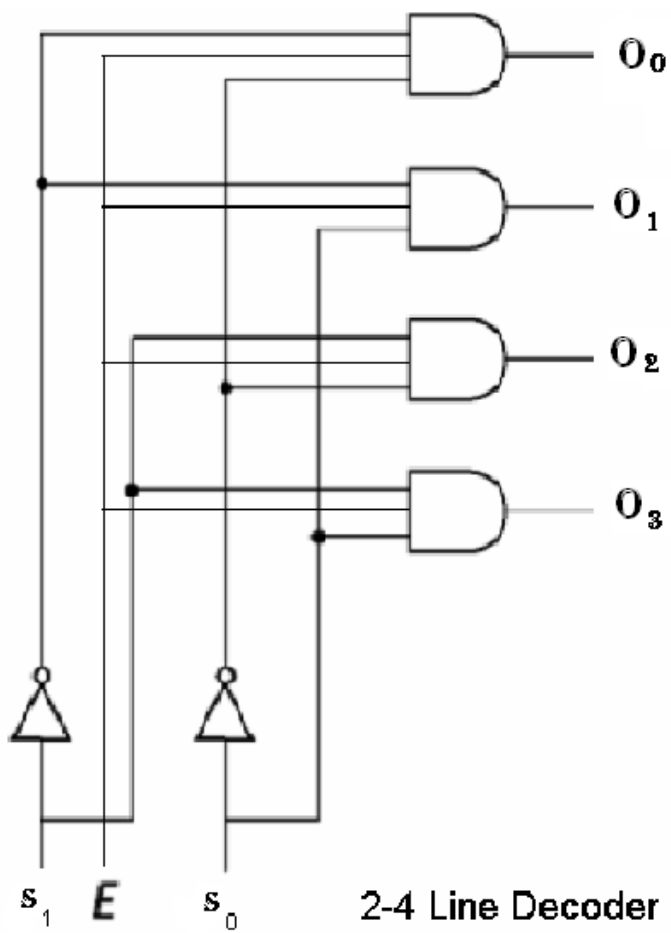
- Many devices have an additional **enable input**, which is used to “**activate**” or “**deactivate**” the device.



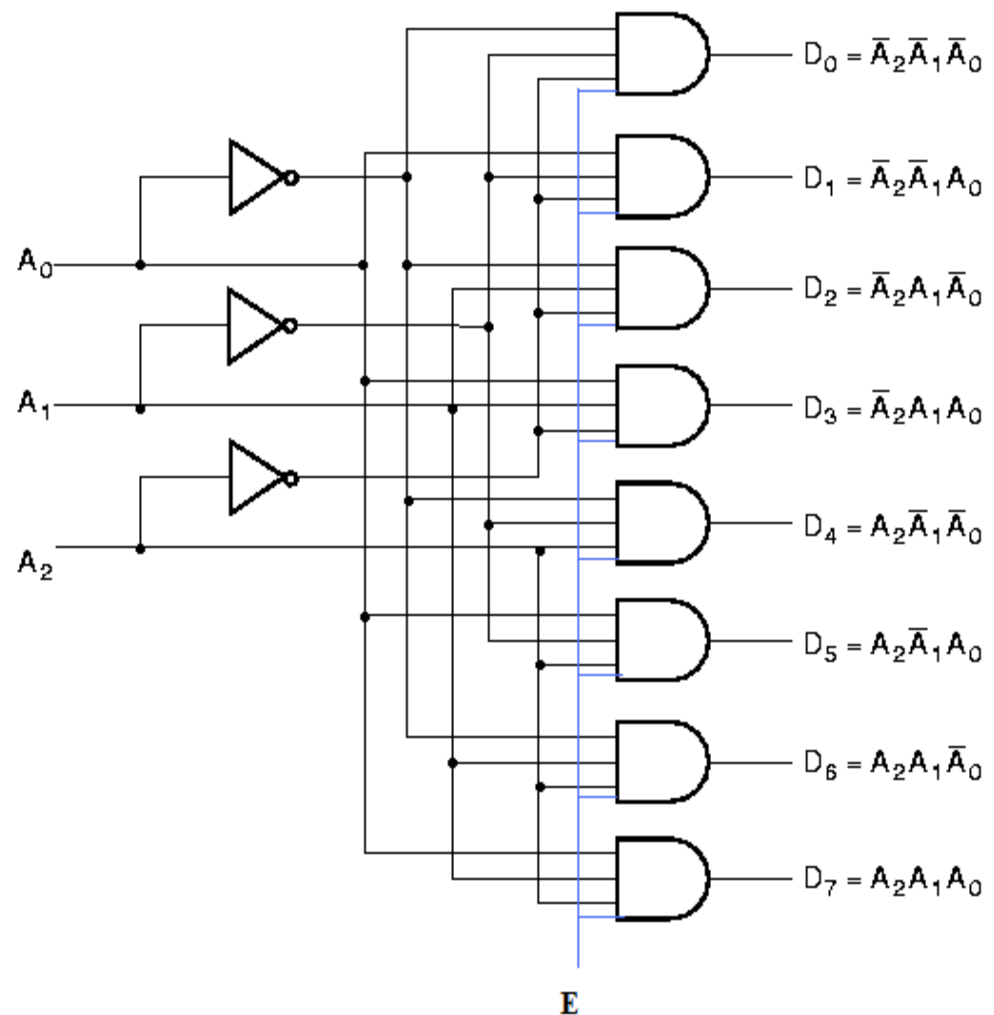
- For a decoder,
 - EN=0 “deactivates” the decoder. By convention, that means *all* of the decoder’s outputs are 0.
 - EN=1 activates the decoder, so it behaves as specified earlier. Exactly one of the outputs will be 1.

EN	S1	S0	Q0	Q1	Q2	Q3
0	0	0	0	0	0	0
0	0	1	0	0	0	0
0	1	0	0	0	0	0
0	1	1	0	0	0	0
1	0	0	1	0	0	0
1	0	1	0	1	0	0
1	1	0	0	0	1	0
1	1	1	0	0	0	1

2-to-4

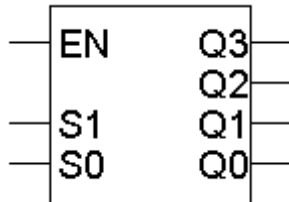


3-to-8



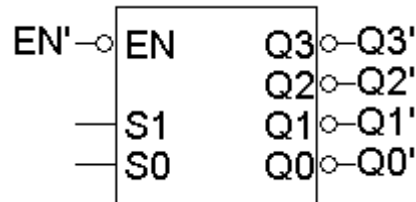
A variation of the standard decoder

- The decoders we've seen so far are **active-high** decoders.



EN	S1	S0	Q0	Q1	Q2	Q3
0	x	x	0	0	0	0
1	0	0	1	0	0	0
1	0	1	0	1	0	0
1	1	0	0	0	1	0
1	1	1	0	0	0	1

- An **active-low decoder** is the same thing, but with an inverted EN input and inverted outputs.



EN	S1	S0	Q0	Q1	Q2	Q3
0	0	0	0	1	1	1
0	0	1	1	0	1	1
0	1	0	1	1	0	1
0	1	1	1	1	1	0
1	x	x	1	1	1	1

So what good is a decoder?

- Do the truth table and equations look familiar?

S1	S0	Q0	Q1	Q2	Q3
0	0	1	0	0	0
0	1	0	1	0	0
1	0	0	0	1	0
1	1	0	0	0	1

$$Q0 = S1' S0'$$

$$Q1 = S1' S0$$

$$Q2 = S1 S0'$$

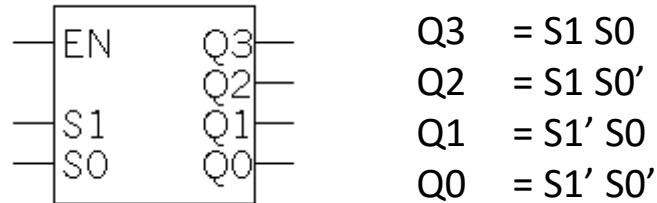
$$Q3 = S1 S0$$

- Decoders are sometimes called **minterm generators**.
 - For each of the input combinations, exactly one output is true.
 - Each output equation contains all of the input variables.
 - These properties hold for all sizes of decoders.
- This means that you can implement arbitrary functions with decoders. **If you have a sum of minterms equation for a function, you can easily use a decoder (a minterm generator) to implement that function.**

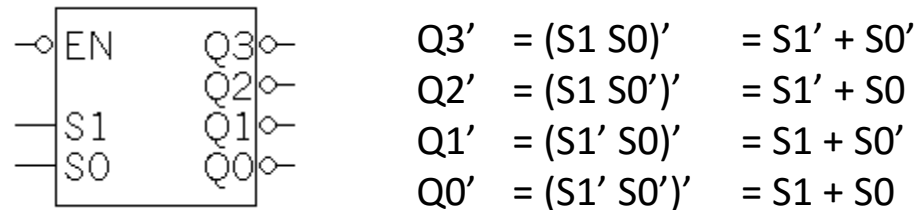


Mysteriously Similar

- **Active-high** decoders generate *minterms*, as we've already seen.



- The output equations for an **active-low** decoder are mysteriously similar, yet somehow different.



- It turns out that **active-low** decoders generate *maxterms*.

Implementing Functions using Decoders

Design example: addition

- Let's make a circuit that adds three 1-bit inputs X, Y and Z.
- We will need two bits to represent the total; let's call them C and S, for "carry" and "sum." Note that C and S are two separate functions of the same inputs X, Y and Z.

$0 + 1 + 1 = 10$



X	Y	Z	C	S
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

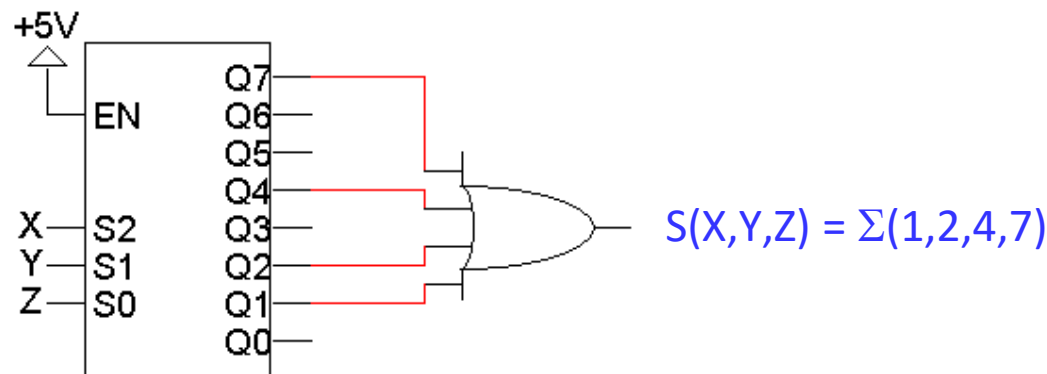
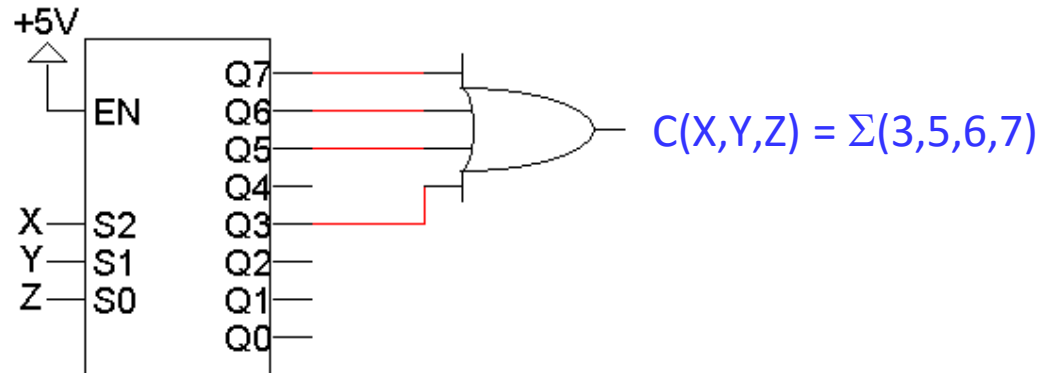
$$C(X,Y,Z) = \Sigma (3,5,6,7)$$

$$S(X,Y,Z) = \Sigma (1,2,4,7)$$

$\leftarrow 1 + 1 + 1 = 11$

Decoder-based adder

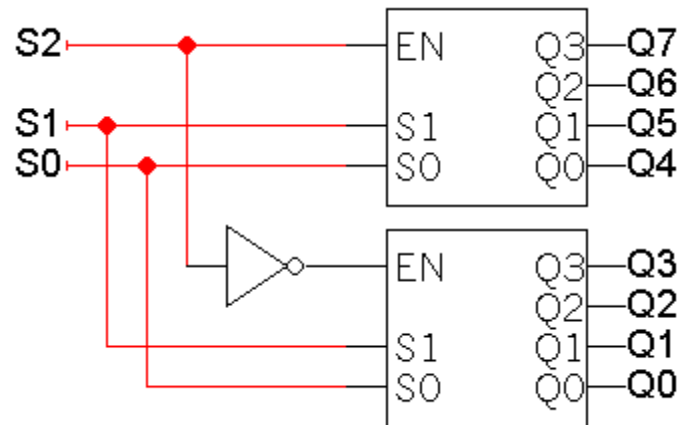
- Here, two 3-to-8 decoders implement C and S as sums of minterms.



- The “+5V” symbol (“5 volts”) is how you represent a constant 1 or true in Logic Works. It has been used here so that the decoders are always active.

Decoder expansion

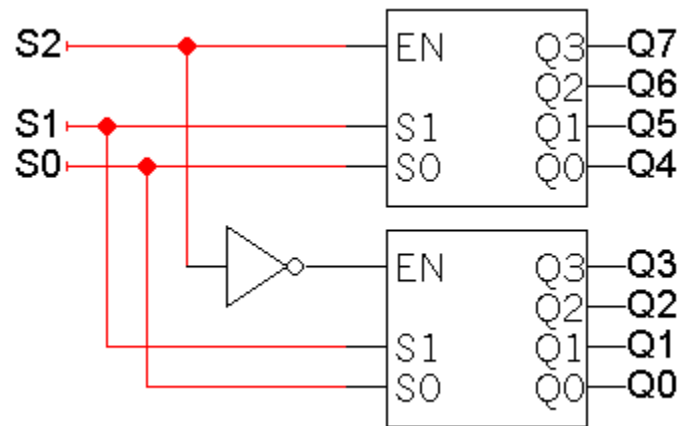
- Combine two or more small decoders with **enable** inputs to form a larger decoder.
- Here a 3-to-8 decoder has been constructed from **two** 2-to-4 decoders:



S2	S1	S0	Q0	Q1	Q2	Q3	Q4	Q5	Q6	Q7
0	0	0	1	0	0	0	0	0	0	0
0	0	1	0	1	0	0	0	0	0	0
0	1	0	0	0	1	0	0	0	0	0
0	1	1	0	0	0	1	0	0	0	0
1	0	0	0	0	0	0	1	0	0	0
1	0	1	0	0	0	0	0	1	0	0
1	1	0	0	0	0	0	0	0	1	0
1	1	1	0	0	0	0	0	0	0	1

Modularity

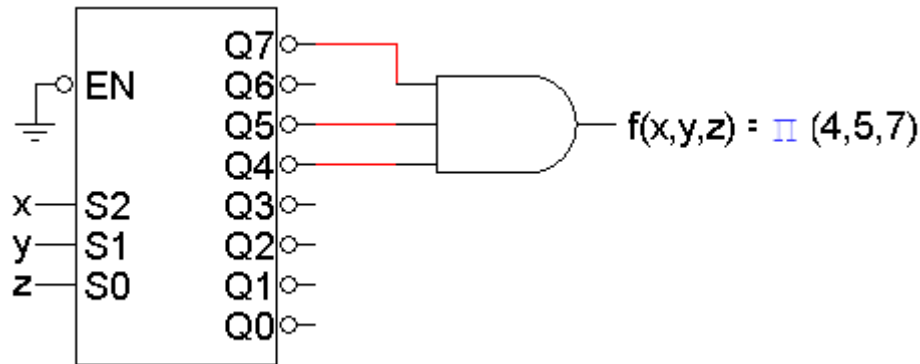
- Be careful not to confuse the “inner” inputs and outputs of the 2-to-4 decoders with the “outer” inputs and outputs of the 3-to-8 decoder (which are in boldface).
- This is similar to having several functions in a program which all use a formal parameter “x”.



- You could verify that this circuit is a 3-to-8 decoder, by using equations for the 2-to-4 decoders to derive equations for the 3-to-8.

Active-low decoder example

- So we can use active-low decoders to implement arbitrary functions too, but as a product of maxterms.
- For example, here is an implementation of the function, $f(x,y,z) = \prod(4,5,7)$, using an active-low decoder.



- The “ground” symbol connected to EN represents logical 0, so this decoder is always enabled.
- Remember that you need an AND gate for a product of sums.

Use two 3 to 8 decoders to make 4 to 16 decoder

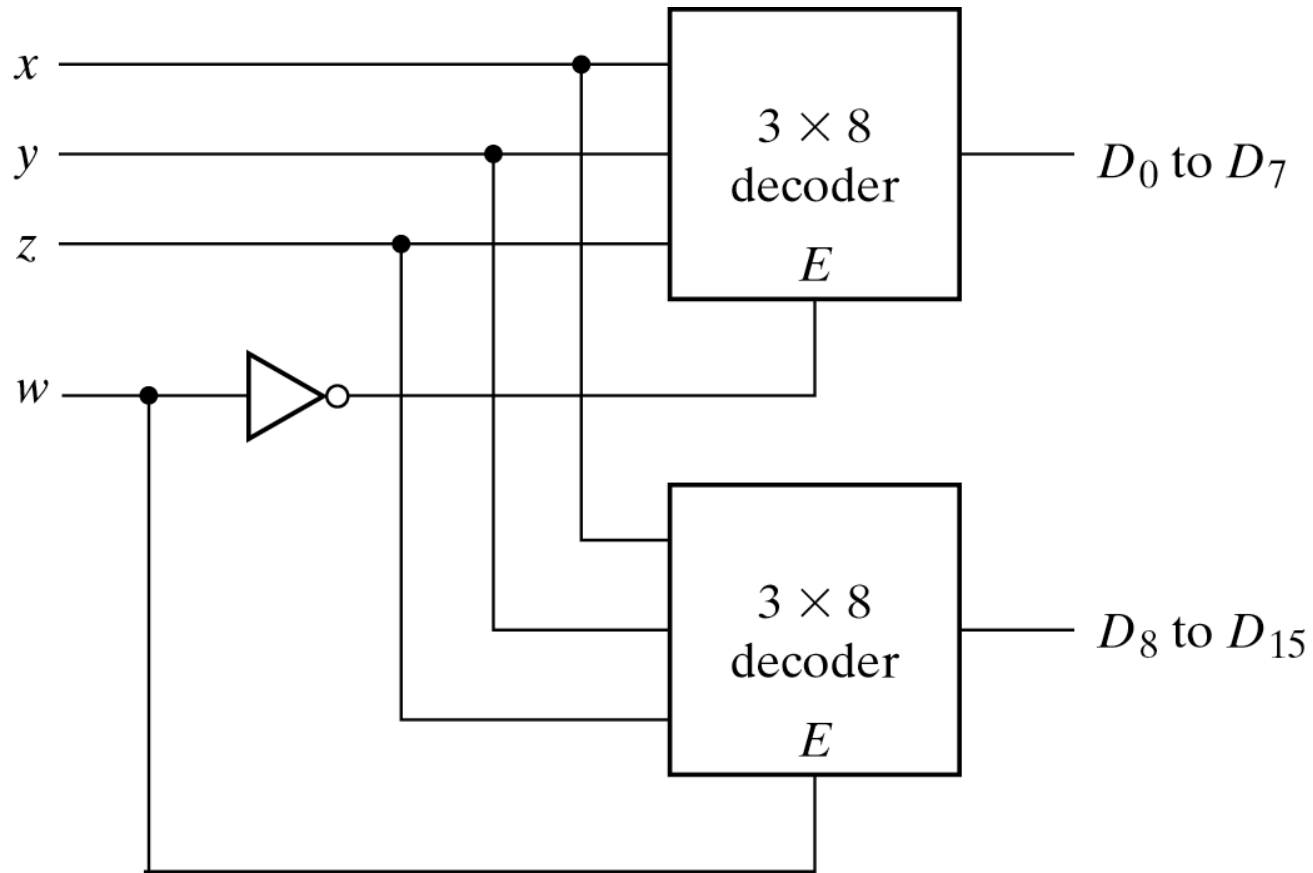
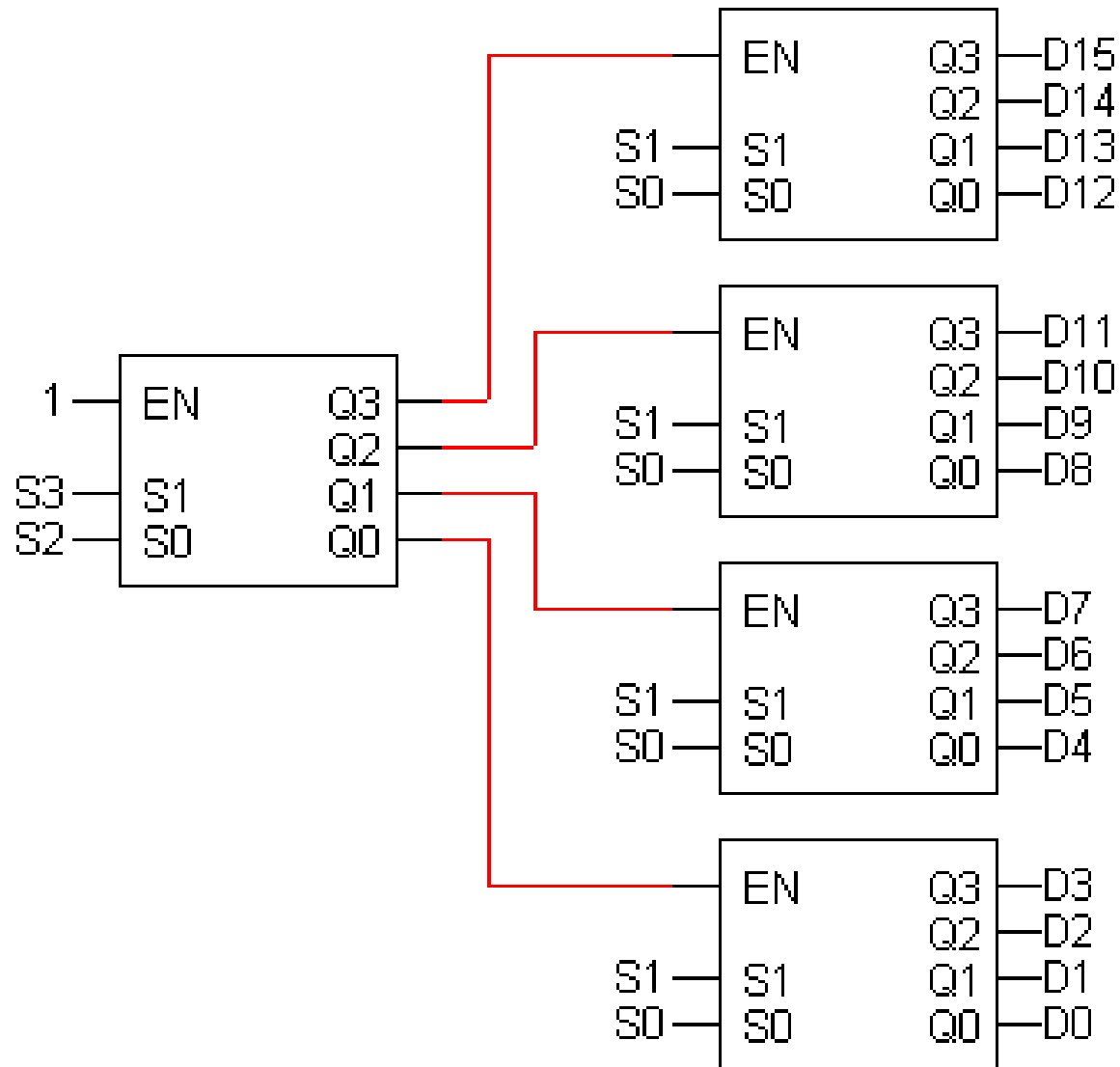
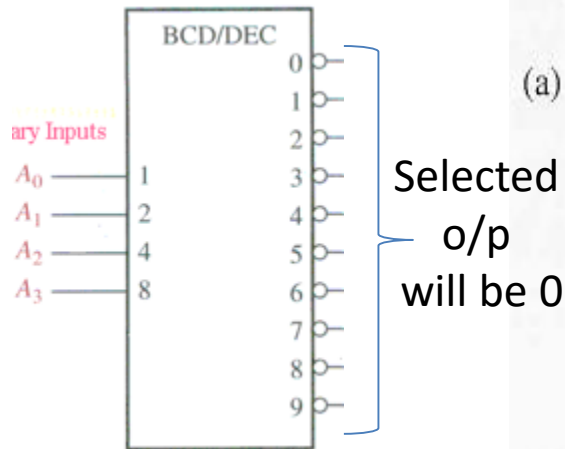


Fig. 4-20 4 × 16 Decoder Constructed with Two 3 × 8 Decoders

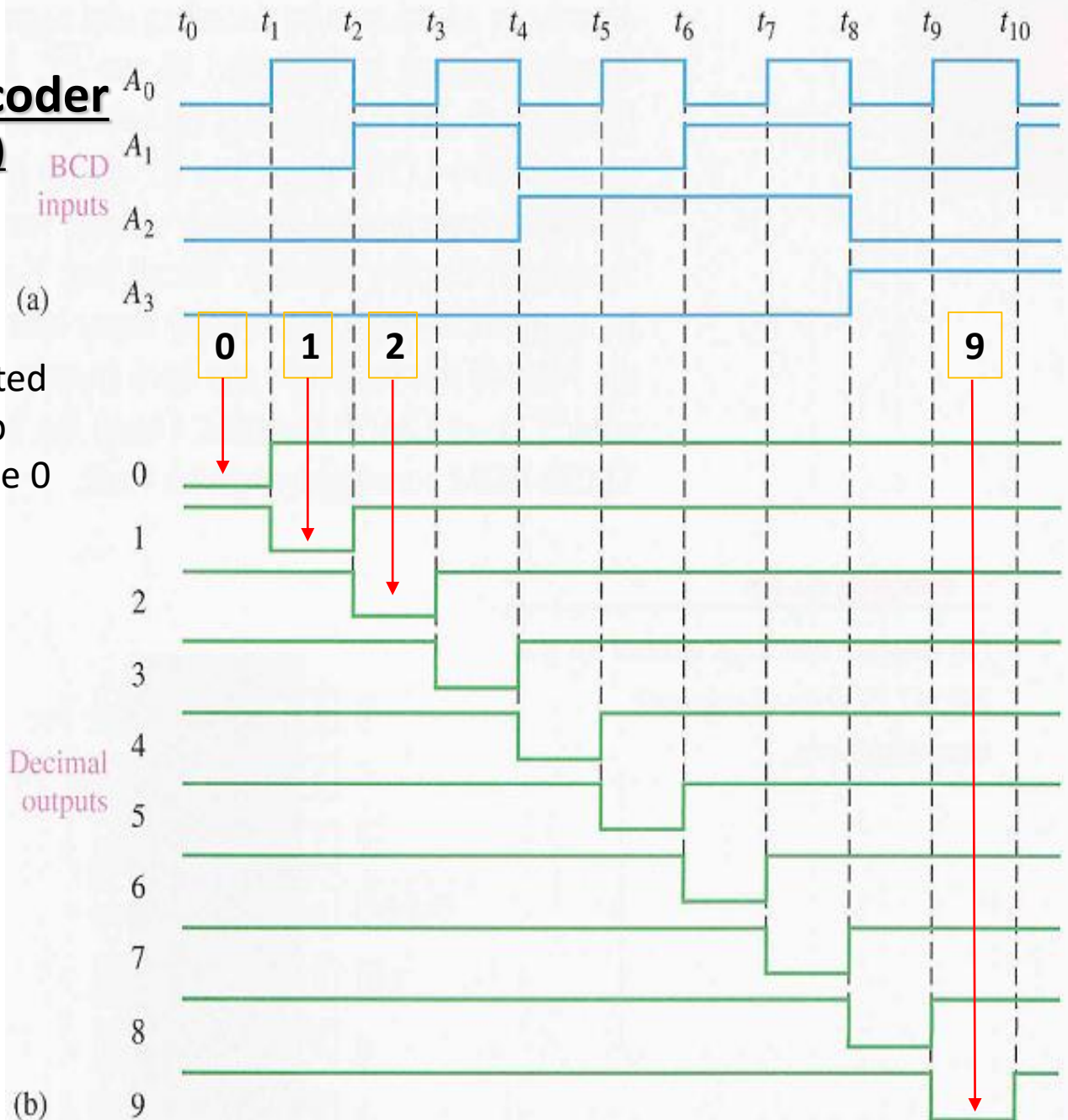
4-to-16 decoder using only 2-to-4 decoders (no gates)



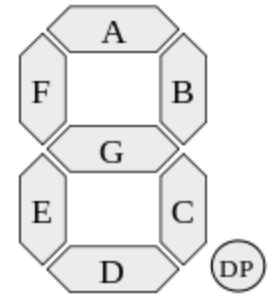
BCD-to-decimal decoder (Active Low Output)



DECIMAL DIGIT	BCD CODE			
	A_3	A_2	A_1	A_0
0	0	0	0	0
1	0	0	0	1
2	0	0	1	0
3	0	0	1	1
4	0	1	0	0
5	0	1	0	1
6	0	1	1	0
7	0	1	1	1
8	1	0	0	0
9	1	0	0	1

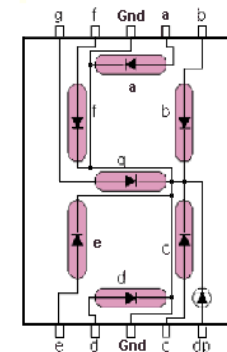


BCD-to-7-segement decoder

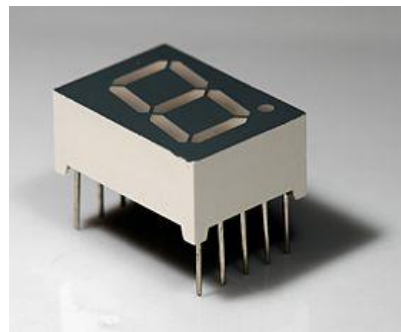
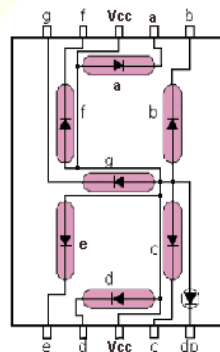


DECIMAL DIGIT	INPUTS				SEGMENT OUTPUTS						
	<i>D</i>	<i>C</i>	<i>B</i>	<i>A</i>	<i>a</i>	<i>b</i>	<i>c</i>	<i>d</i>	<i>e</i>	<i>f</i>	<i>g</i>
0	0	0	0	0	1	1	1	1	1	1	0
1	0	0	0	1	0	1	1	0	0	0	0
2	0	0	1	0	1	1	0	1	1	0	1
3	0	0	1	1	1	1	1	1	0	0	1
4	0	1	0	0	0	1	1	0	0	1	1
5	0	1	0	1	1	0	1	1	0	1	1
6	0	1	1	0	1	0	1	1	1	1	1
7	0	1	1	1	1	1	1	0	0	0	0
8	1	0	0	0	1	1	1	1	1	1	1
9	1	0	0	1	1	1	1	1	0	1	1
10	1	0	1	0	X	X	X	X	X	X	X
11	1	0	1	1	X	X	X	X	X	X	X
12	1	1	0	0	X	X	X	X	X	X	X
13	1	1	0	1	X	X	X	X	X	X	X
14	1	1	1	0	X	X	X	X	X	X	X
15	1	1	1	1	X	X	X	X	X	X	X

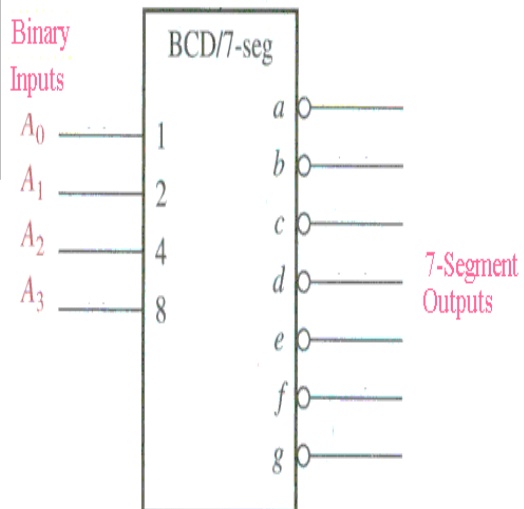
Common Cathode

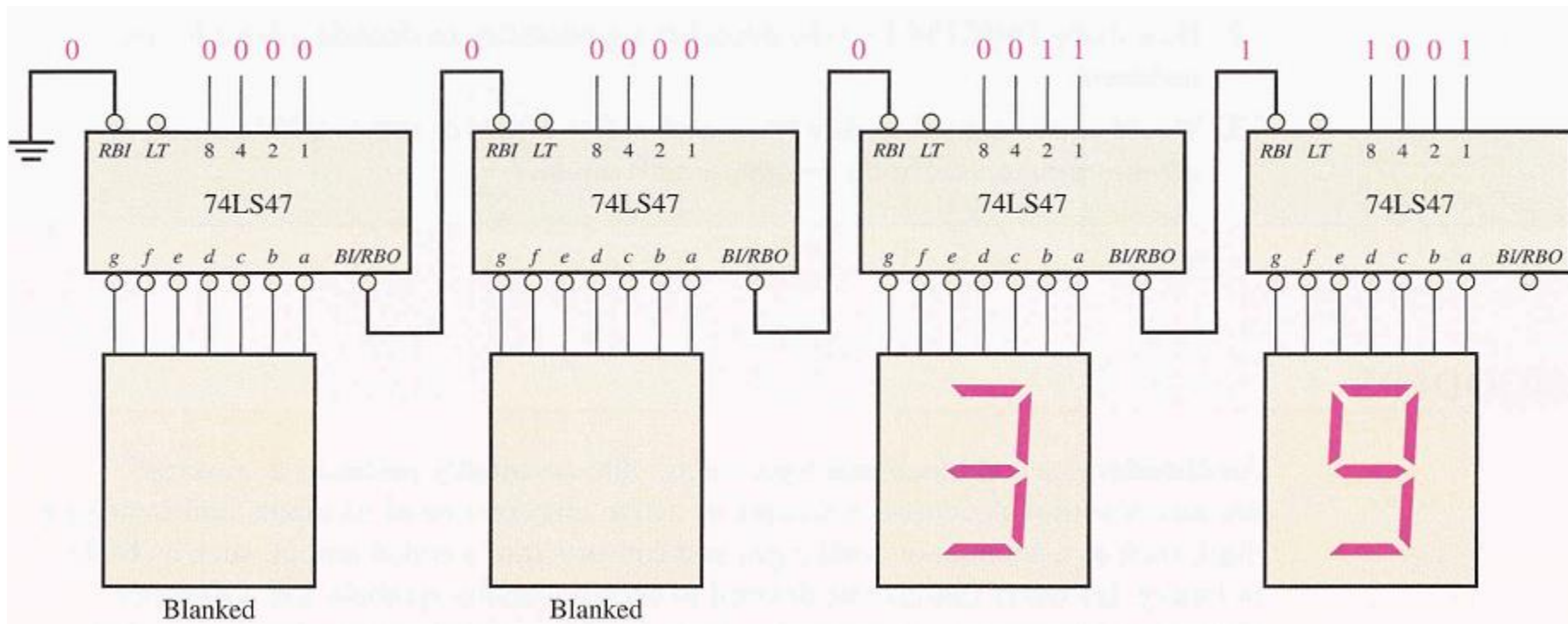


Common Anode



Binary Inputs







Summary

- A n -to- 2^n decoder generates the minterms of a n -variable function.
 - As such, decoders can be used to implement arbitrary functions.
- Some variations of the basic decoder include:
 - Adding an enable input.
 - Using active-low inputs and outputs to generate maxterms.
- We also talked about:
 - Applying our circuit analysis and design techniques to understand and work with decoders.
 - Using block symbols to encapsulate common circuits like decoders.
 - Building larger decoders from smaller ones.

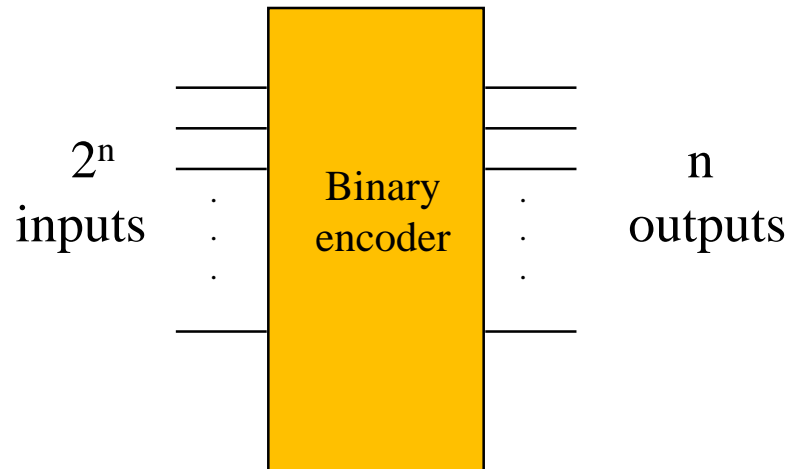
Encoders

An **Encoder** is a combinational logic circuit that performs a “**reverse**” decoder function.

An Encoder accepts an active level on one of its inputs representing a digit, such as a decimal or octal digit, and converts it to a coded output, such as BCD or binary.

Encoders can also be devised to encode various symbols and alphabetic characters. The process of converting from familiar symbols or numbers to a coded format is called **Encoding**.

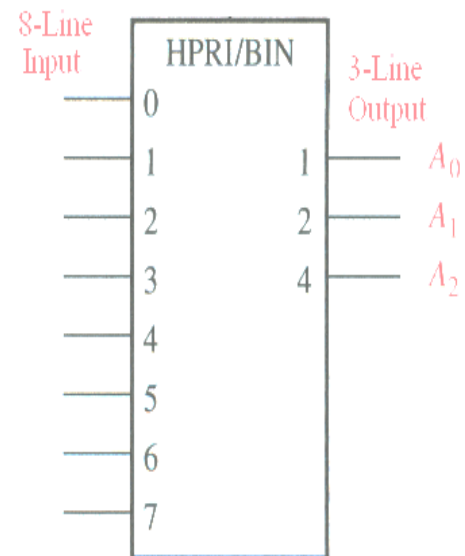
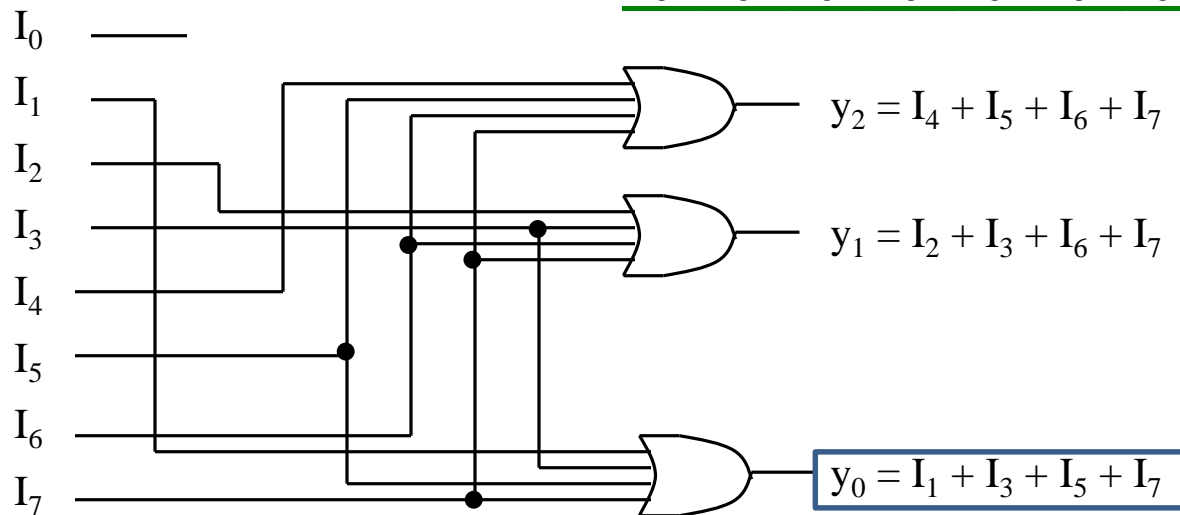
2ⁿ-to-n Encoder:



8-to-3 Binary Encoder

At any one time, only one input line has a value of 1.

Inputs								Outputs		
I_0	I_1	I_2	I_3	I_4	I_5	I_6	I_7	y_2	y_1	y_0
1	0	0	0	0	0	0	0	0	0	0
0	1	0	0	0	0	0	0	0	1	0
0	0	1	0	0	0	0	0	0	1	0
0	0	0	1	0	0	0	0	0	1	0
0	0	0	0	1	0	0	0	1	0	0
0	0	0	0	0	1	0	0	1	0	0
0	0	0	0	0	0	1	0	1	1	0
0	0	0	0	0	0	0	1	1	1	0



8-to-3 Priority Encoder

- What if more than one input line has a value of 1?

Example:

- For the above mentioned problem, let's give priority to higher bits
- Ignore "lower priority" inputs.
- The sequence is:

$$7 > 6 > 5 > 4 > 3 > 2 > 1 > 0$$

- **Idle** indicates that no input is a 1.

Inputs								Outputs			
I ₀	I ₁	I ₂	I ₃	I ₄	I ₅	I ₆	I ₇	Y ₂	Y ₁	Y ₀	Idle
0	0	0	0	0	0	0	0	x	x	x	1
1	0	0	0	0	0	0	0	0	0	0	0
X	1	0	0	0	0	0	0	0	0	1	0
X	X	1	0	0	0	0	0	0	1	0	0
X	X	X	1	0	0	0	0	0	1	1	0
X	X	X	X	1	0	0	0	1	0	0	0
X	X	X	X	X	1	0	0	1	0	1	0
X	X	X	X	X	X	1	0	1	1	0	0
X	X	X	X	X	X	X	1	1	1	1	0

Priority Encoder (8 to 3 encoder)

- Assign priorities to the inputs
- When more than one inputs are asserted, the output generates the code of the input with the highest priority: $7 > 6 > 5 > 4 > 3 > 2 > 1 > 0$

Priority Encoder :

$H7 = I7$ (Highest Priority)

$H6 = I6 \cdot I7'$

$H5 = I5 \cdot I6' \cdot I7'$

$H4 = I4 \cdot I5' \cdot I6' \cdot I7'$

$H3 = I3 \cdot I4' \cdot I5' \cdot I6' \cdot I7'$

$H2 = I2 \cdot I3' \cdot I4' \cdot I5' \cdot I6' \cdot I7'$

$H1 = I1 \cdot I2' \cdot I3' \cdot I4' \cdot I5' \cdot I6' \cdot I7'$

$H0 = I0 \cdot I1' \cdot I2' \cdot I3' \cdot I4' \cdot I5' \cdot I6' \cdot I7'$

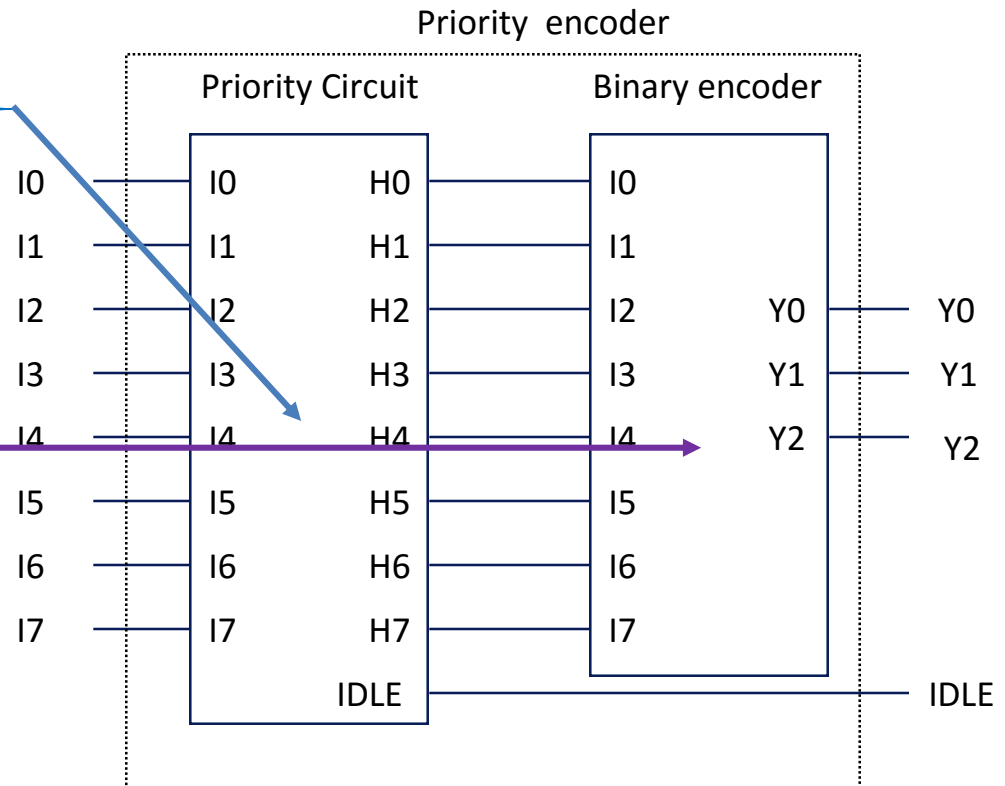
$IDLE = I0' \cdot I1' \cdot I2' \cdot I3' \cdot I4' \cdot I5' \cdot I6' \cdot I7'$

Encoder

$Y0 = I1 + I3 + I5 + I7$

$Y1 = I2 + I3 + I6 + I7$

$Y2 = I4 + I5 + I6 + I7$



Priority Encoder (Irregular sequence)

- Assign priorities to the inputs in the following order:

$2 > 5 > 3 > 4 > 6 > 7 > 1 > 0$

- Priority Encoder :

$H2 = I2$ (Highest Priority)

$H5 = I5 \cdot I2'$

$H3 = I3 \cdot I5' \cdot I2'$

$H4 = I4 \cdot I3' \cdot I5' \cdot I2'$

$H6 = I6 \cdot I4' \cdot I3' \cdot I5' \cdot I2'$

$H7 = I7 \cdot I6' \cdot I4' \cdot I3' \cdot I5' \cdot I2'$

$H1 = I1 \cdot I7' \cdot I6' \cdot I4' \cdot I3' \cdot I5' \cdot I2'$

$H0 = I0 \cdot I1' \cdot I7' \cdot I6' \cdot I4' \cdot I3' \cdot I5' \cdot I2'$

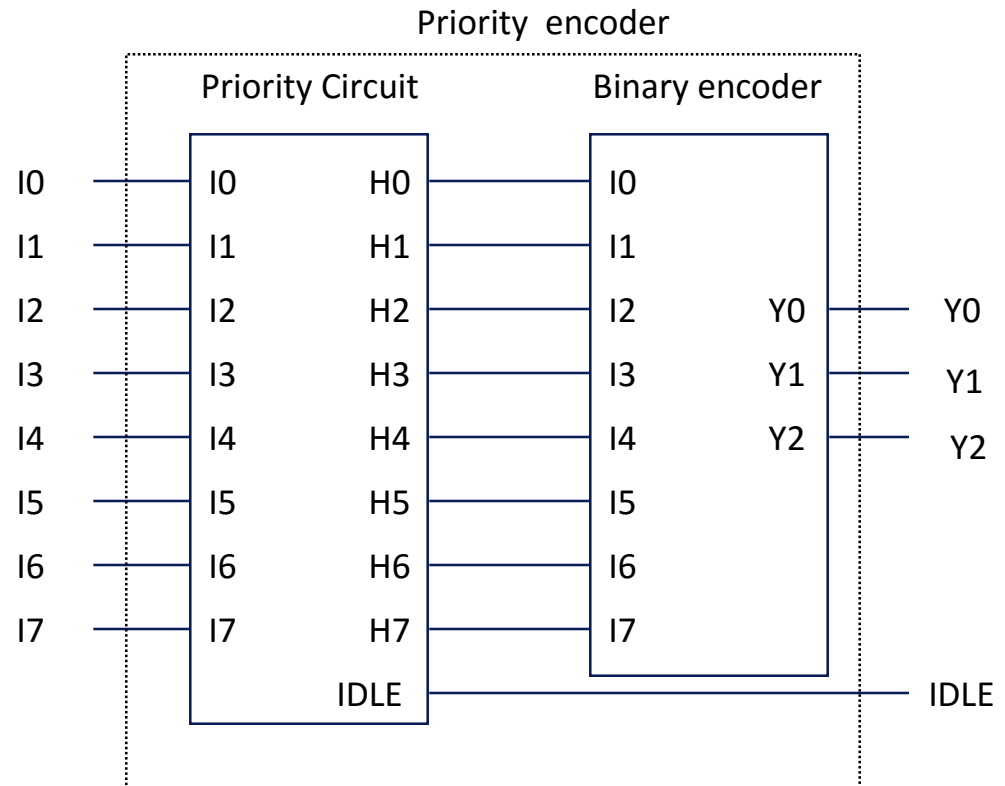
$IDLE = I0' \cdot I1' \cdot I2' \cdot I3' \cdot I4' \cdot I5' \cdot I6' \cdot I7'$

- Encoder

$Y0 = I1 + I3 + I5 + I7$

$Y1 = I2 + I3 + I6 + I7$

$Y2 = I4 + I5 + I6 + I7$



The Decimal-to-BCD Priority Encoder:

Let Priority is given to the higher order digits. Requirements to activate A0:

- 1) A_0 is HIGH if 1 is HIGH and 2,4,6,8 LOW
 A_0 is HIGH if 3 is HIGH and 4,6,8 LOW
 A_0 is HIGH if 5 is HIGH and 6,8 LOW
 A_0 is HIGH if 7 is HIGH and 8 LOW
 A_0 is HIGH if 9 is HIGH

Therefore, $A_0 = 1.2'.4'.6'.8' + 3.4'.6'.8' + 5.6'.8' + 7.8' + 9$

- 2) A_1 is HIGH if 2 is HIGH and 4,5,8,9 LOW
 A_1 is HIGH if 3 is HIGH and 4,5,8,9 LOW
 A_1 is HIGH if 6 is HIGH and 8,9 LOW
 A_1 is HIGH if 7 is HIGH and 8,9 LOW

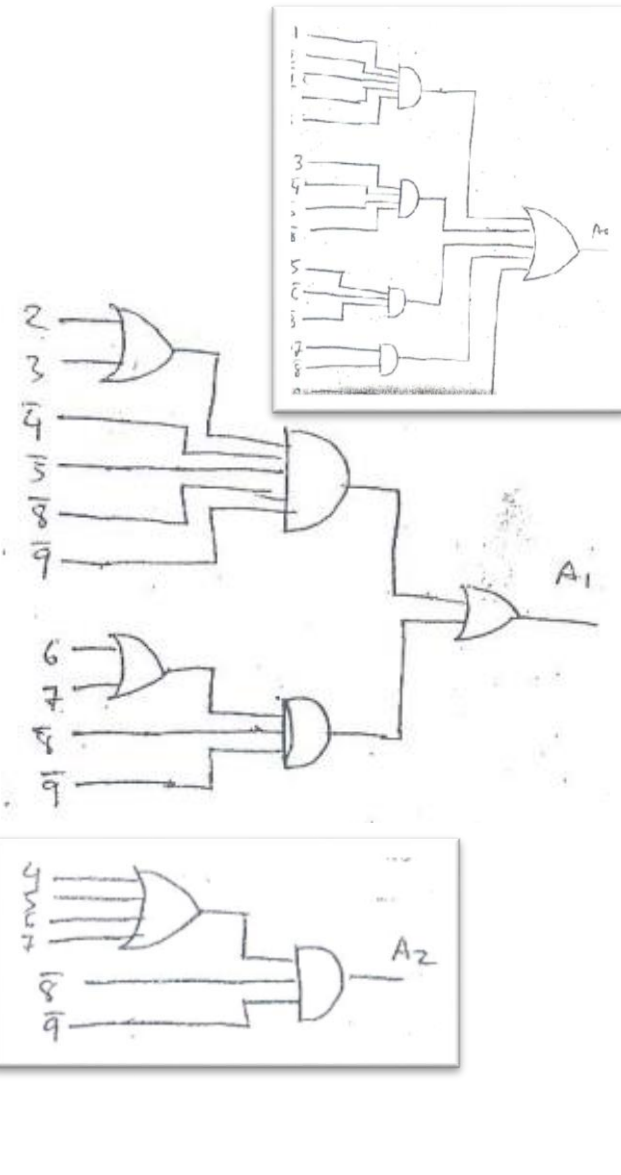
Therefore, $A_1 = (2+3)4'.5'.8'.9' + (6+7)8'.9'$

- 3) A_2 is HIGH if 4 is HIGH and 8,9 LOW
 A_2 is HIGH if 5 is HIGH and 8,9 LOW
 A_2 is HIGH if 6 is HIGH and 8,9 LOW
 A_2 is HIGH if 7 is HIGH and 8,9 LOW

Therefore, $A_2 = (4+5+6+7)8'. 9'$

- 4) A_3 is HIGH if 8 & 9 are HIGH

Therefore, $A_3 = 8+9$





Reference:

**Mixed contents from books by Floyd; Mano; Vahid
And Howard.**

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Thanks