## Representing Graphs and Graph Isomorphism



Course Code: 00090 Course Title: Discrete Mathematics

Dept. of Computer Science Faculty of Science and Technology

Lecturer No:	16	Week No:	9	Semester:	Summer 21-22	
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## Lecture Outline



Representing Graphs and Graph Isomorphism (8.3)

- Graph Representation:
  - Adjacency lists
  - Adjacency matrices
  - Incidence matrices
- Graph Isomorphism

## Objectives and Outcomes



 Objectives: To understand three different representation of a Graph, to understand graph invariants and graph isomorphism.

 Outcomes: This students are expected to be able to represent a graph using adjacency list, adjacency matrix, and incidence matrix; be able to determine whether two graphs are isomorphic.

## Representing Graphs



- One way to represent a graph without multiple edges is to list all the edges of the graph.
- Another way to represent a graph with no multiple edges is to use adjacency lists, which specify the vertices that are adjacent to each vertex of the graph.
- Adjacency Lists: A table with 1 row per vertex, listing its adjacent vertices.
- <u>Directed Adjacency Lists</u>: A table with 1 row per node, listing the terminal nodes of each edge incident from that node.

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## Representing a Graph with an Adjacency List

• EXAMPLE 1: Use adjacency lists to describe the simple graph given in Figure 1.

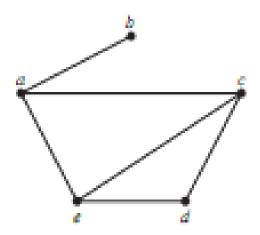


FIGURE 1 A Simple Graph.

TABLE 1 An Adjacency List for a Simple Graph.				
Vertex	Adjacent Vertices			
а	b, c, e			
b	а			
c	a, d, e			
d	€, €			
e	a, c, d			



## Representing a Graph with an Adjacency List

EXAMPLE 2: Represent the directed graph shown in Figure 2
by listing all the vertices that are the terminal vertices of
edges starting at each vertex of the graph.

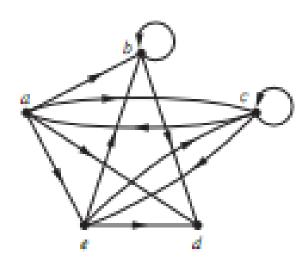


FIGURE 2 A Directed Graph.

TABLE 2 An Adjacency List for a Directed Graph.					
Initial Vertex	Terminal Vertices				
а	b, c, d, e				
ь	b,d				
¢	a,c,e				
d					
e	b, c, d				



## **Representing Graphs using Matrices**

- Two types of matrices commonly used to represent graphs
  - 1) Adjacency matrix
  - 2) Incidence matrix
- Adjacency matrix: A matrix representing a graph using the adjacency of vertices.
- Incidence matrix: A matrix representing a graph using the incidence of edges and vertices.



- Suppose that G = (V, E) is a simple graph where |V| = n.
   Suppose that the vertices of G are listed arbitrarily as v<sub>1</sub>, v<sub>2</sub>, ...., v<sub>n</sub>.
  - The **adjacency matrix A** of G, with respect to this listing of the vertices, is the  $n \times n$  zero-one matrix with 1 as its (i,j)th entry when  $v_i$  and  $v_j$  are adjacent, and 0 as its (i,j)th entry when they are not adjacent.
- In other words, if its adjacency matrix is  $A = [a_{ij}]$ , then  $a_{ij} = [1]$  if  $\{v_i, v_j\}$  is an edge of G,
  - 0 otherwise



#### **Example 3: Representing a graph using an Adjacency Matrix**

EXAMPLE 3 Use an adjacency matrix to represent the graph shown in Figure 3.

Solution: We order the vertices as a, b, c, d. The matrix representing this graph is

$$\begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{bmatrix}.$$

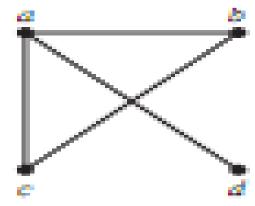


FIGURE 3

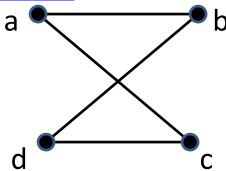


## Given an Adjacency matrix, draw the graph

• Example 4: Draw a graph with the adjacency matrix below.

Γο	1	1	0 1 1
0 1 1 0	0	0	1
1	0	0	1
Lo	1	1	o _

### **Solution**:

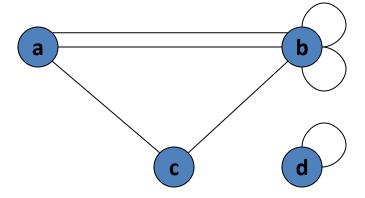




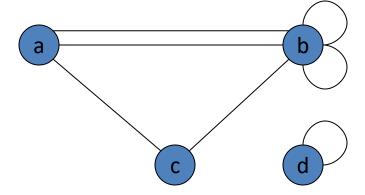
- Adjacency matrix can also be used to represent undirected graphs with loops and with multiple edges.
  - A loop at the vertex  $a_i$  is represented by 1 at the (i,j)th position of the adjacency matrix.
  - When multiple edges are present, the adjacency matrix is no longer zero-one matrix, because the (i,j)th entry of this matrix equals the number of edges that are associated to {a<sub>i</sub>,a<sub>j</sub>}
  - All undirected graphs, including multigraphs and pseudographs, have symmetric adjacency matrices



<u>Example</u>: What's the adjacency matrix of the following graph?







## **Solution:**

$$egin{pmatrix} 0 & 2 & 1 & 0 \ 2 & 2 & 1 & 0 \ 1 & 1 & 0 & 0 \ 0 & 0 & 0 & 1 \end{pmatrix}$$



### **Practice at Home**

• **Example 5**: Use an adjacency matrix to represent the pseudograph shown in **Figure 5**.

## Adjacency matrices for Directed Graphs

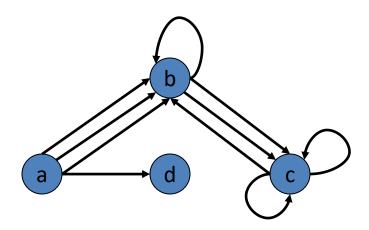


- Adjacency matrices can also be used to represent directed multigraphs.
  - However, such matrices are not zero-one matrices when there are multiple edges in the same direction connecting two vertices
- In the adjacency matrix for a directed graph,  $a_{ij}$  equals the number of edges that are associated to  $(v_{i}, v_{i})$ .

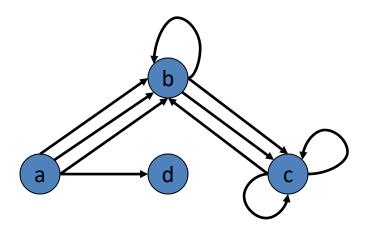


## Adjacency matrices for directed graphs

**Example:** What is the adjacency matrix?







## **Solution:**

$$\begin{pmatrix}
0 & 3 & 0 & 1 \\
0 & 1 & 2 & 0 \\
0 & 1 & 2 & 0 \\
0 & 0 & 0 & 0
\end{pmatrix}$$

## Incidence matrices



- Incidence matrix: A matrix representing a graph using the incidence of edges and vertices.
- Let G = (V, E) be an undirected graph. Suppose that  $v_1$ ,  $v_2$ , ....,  $v_n$  are the vertices and  $e_1$ ,  $e_2$ , ....,  $e_m$  are the edges of G.

Then the incidence matrix with respect to this ordering of V and E is the  $n \times m$  matrix,  $M = [m_{ii}]$ , where

$$m_{ij} = \begin{bmatrix} 1 & \text{when edge } e_j \text{ is incident with } v_i, \\ 0 & \text{otherwise} \end{bmatrix}$$

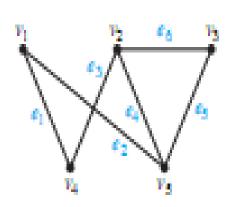


### **Example 6: Representing a graph with an Incidence Matrix**

$$m_{lj} = \begin{cases} 1 & \text{when edge } e_j \text{ is incident with } v_l, \\ 0 & \text{otherwise.} \end{cases}$$

EXAMPLE 6 Represent the graph shown in Figure 6 with an incidence matrix.

Solution: The incidence matrix is





### **Example 7: Representing a graph with an Incidence**

## EXAMPLE 7 Represent the pseudograph shown in Figure 7 using an incidence matrix.



→" Solution: The incidence matrix for this graph is

## FIGURE 7

A Pseudograph.

## Isomorphism of Graphs



• <u>Definition 1</u>: The simple graphs  $G_1 = (V_1, E_1)$  and  $G_2 = (V_2, E_2)$  are isomorphic **if there is a one-to-one and onto function f from V\_1 to V\_2** with the property that a and b are adjacent in  $G_1$  iff f(a) and f(b) are adjacent in  $G_1$ , for all a and b in  $V_1$ .

#### Such a function **f** is called an **isomorphism**.

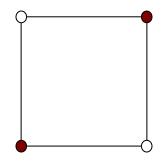
- When two simple graphs are isomorphic, there is a one-to-one correspondence between vertices of the two graphs that preserves the adjacency relationship
- Two graphs are isomorphic iff they are identical except for their node names
- Isomorphism of simple graphs is an equivalence relation

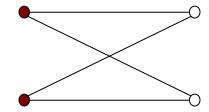


## **Isomorphism of Graphs**

• Intuitively, two graphs are isomorphic if can bend, stretch and reposition vertices of the first graph, until the second graph is formed. Etymologically, isomorphic means "same shape".

e.g.: Can twist or relabel:



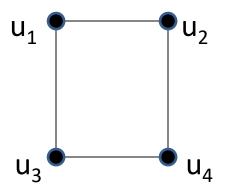


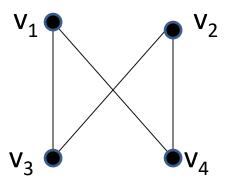
to obtain:



## **Example 8**

• Show that the graphs G=(V,E) and H=(W,F) are isomorphic.





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## **Solution of Example 8**

- The function f with  $f(u_1) = v_1$ ,  $f(u_2) = v_4$ ,  $f(u_3) = v_3$ , and  $f(u_4) = v_2$  is a one-to-one correspondence between V and W. To see that this correspondence preserves adjacency, note that adjacent vertices in G are  $u_1$  and  $u_2$ ,  $u_1$  and  $u_3$ ,  $u_2$  and  $u_4$ , and each of the pairs  $f(u_1) = v_1$  and  $f(u_2) = v_4$ ,  $f(u_1) = v_1$  and  $f(u_3) = v_3$ ,  $f(u_2) = v_4$  and  $f(u_4) = v_2$ , and  $f(u_3) = v_3$  and  $f(u_4) = v_2$  are adjacent in H.
- So, the graphs G and H are isomorphic.

## Isomorphism of Graphs

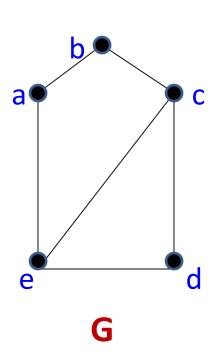


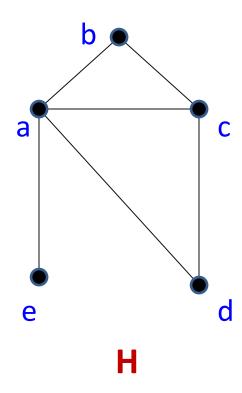
- It is often difficult to determine whether two simple graphs are isomorphic.
- Sometimes it is not hard to show that two graphs are not isomorphic.
- A property preserved by isomorphism of graph is called a graph invariant.
- Isomorphic simple graphs must have
  - The same # of vertices
  - The same # of edges
  - The degrees of the vertices must be the same and degrees of adjacent vertices must be same
  - Same length of simple circuit(s), if any



## **Example 9**

Show that the following two graphs G and H are not isomorphic.







## **Solution of Example 9**

 Solution: Both G and H have five vertices and six edges. However, H has a vertex of degree one, namely e, whereas G has no vertices of degree one. It follows that G and H are not isomorphic.

There are other answers too!
 What are the other answers?

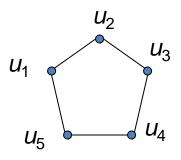


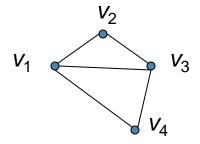
## **Isomorphism of Graphs**

- Note: The number of vertices, the number of edges, and the number of vertices of each degree are all invariants under isomorphism. If any of these quantities differ in two simple graphs, these graphs cannot be isomorphic.
- However, when these invariants are the same, it does not necessarily mean that the two graphs are isomorphic. There are NO useful sets of invariants currently known that can be used to determine whether simple graphs are isomorphic.



Q1: Why are the following graphs not isomorphic?

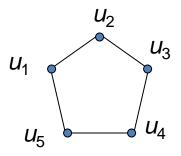


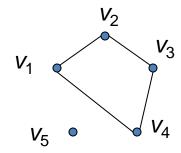




A1: 1<sup>st</sup> graph has more vertices than 2<sup>nd</sup>. So the graphs are not isomorphic. What are other answers?

Q2: Why are the following graphs not isomorphic?

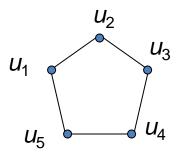


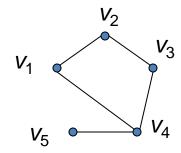




A2: 1<sup>st</sup> graph has more edges than 2<sup>nd</sup>. What are other answers?

Q3: Why are the following graphs not isomorphic?

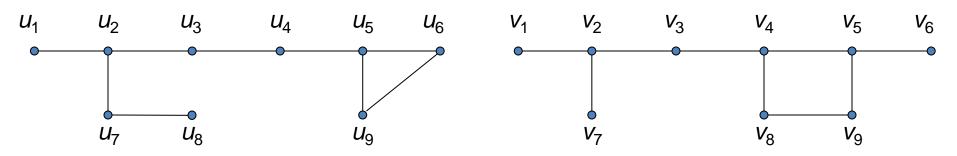






A3: 2<sup>nd</sup> graph has vertex of degree 1, 1<sup>st</sup> graph doesn't. So the graphs are not isomorphic. What are other answers?

Q4: Why are the following graphs non-isomorphic?





• A4: 1<sup>st</sup> graph has 2 vertices of degree 1, whereas the 2<sup>nd</sup> graph has 3 vertices of degree 1.

So the graphs are not isomorphic.

What are other answers?

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## **Practice @ Home**

- Example 10
- Example 11
- Relevant Exercises from your text book

#### **Books**



- Rosen, K. H., & Krithivasan, K. (2012). Discrete mathematics and its applications: with combinatorics and graph theory. Tata McGraw-Hill Education. (7<sup>th</sup> Edition)
- Liu, C. L. (1986). Elements of discrete mathematics. Tata McGraw-Hill Education.

#### References



- 1. Discrete Mathematics, *Richard Johnsonbaugh*, Pearson education, Inc.
- 2. Discrete Mathematical Structures, *Bernard Kolman*, *Robert C. Busby*, *Sharon Ross*, Prentice-Hall, Inc.
- 3. SCHAUM'S outlines Discrete Mathematics(2<sup>nd</sup> edition), by Seymour Lipschutz, Marc Lipson
- An online tutorial <a href="https://www.geeksforgeeks.org/graph-and-its-representations/">https://www.geeksforgeeks.org/graph-and-its-representations/</a>
- Another online tutorial <a href="https://www.gatevidyalay.com/graph-isomorphism/">https://www.gatevidyalay.com/graph-isomorphism/</a>
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