# CSC 2221: Algorithms

Lecture 3

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- Bubble Sort
- 2 Insertion sort
- Selection sort
- 4 Linear Search



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## **Bubble Sort**

- Running time  $O(n^2)$
- a[i] > a[i+1] swap(a[i], a[i+1])
- in the loop, index i=0 to ?

20	9	6	3	1
9	20	6	3	1
9	6	20	3	1
9	6	3	20	1
9	6	3	1	20



# Bubble Sort Algorithm

#### Algorithm 5 Bubble Sort

```
1: procedure BubbleSort(A, n)
      for k \leftarrow 0, n-1 do
          for i \leftarrow 0, n-2 do
3:
             if A[i] > A[i+1] then
4:
                swap(A[i], A[i+1]
5:
             end if
6:
          end for
7:
      end for
9: end procedure
```



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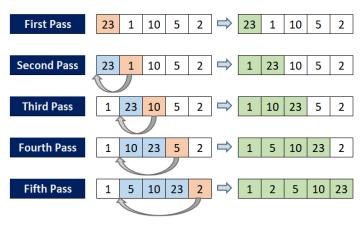
### Insertion sort

#### Algorithm 7 Insertion Sort

```
1: procedure InsertionSort(A, n)
       for i \leftarrow 1, n-1 do
          value \leftarrow A[i]
 3:
       i \leftarrow j-1
 4:
           while i \ge 0 \& A[i] > value do
5:
               swap(A[i], A[i+1])
6:
              i \leftarrow i - 1
7:
           end while
           A[i+1] = value
9:
       end for
10:
11: end procedure
```



### Insertion sort simulation





# Insertion sort complexity

```
INSERTION-SORT (A)
                                                     times
                                             cost
   for j = 2 to A. length
                                             C_1
2 kev = A[i]
                                                     n-1
                                             C_2
      // Insert A[j] into the sorted
           sequence A[1 ... j - 1].
                                                     n-1
                                             c_4 n-1
      i = i - 1
                                                    \sum_{j=2}^{n} t_j
      while i > 0 and A[i] > key
                                             C_5
                                             c_6 \qquad \sum_{j=2}^n (t_j - 1)
6
           A[i+1] = A[i]
                                                  \sum_{j=2}^{n} (t_j - 1)
           i = i - 1
                                             C_7
8
      A[i+1] = key
                                             C_{8}
```

## Insertion sort complexity

$$T(n) = c_1 n + c_2(n-1) + c_4(n-1) + c_5 \sum_{j=2}^{n} t_j + c_6 \sum_{j=2}^{n} (t_j - 1) + c_7 \sum_{j=2}^{n} (t_j - 1) + c_8(n-1)$$

Best Case : List is sorted

$$T(n) = c_1 n + c_2 (n-1) + c_4 (n-1) + c_5 \sum_{j=2}^{n} 1 + c_8 (n-1)$$

$$= c_1 n + c_2 (n-1) + c_4 (n-1) + c_5 (n-1) + c_8 (n-1)$$

$$= (c_1 + c_2 + c_4 + c_5 + c_8) n - (c_2 + c_4 + c_5 + c_8)$$

$$= an + b$$



## Insertion sort complexity

$$T(n) = c_1 n + c_2(n-1) + c_4(n-1) + c_5 \sum_{j=2}^{n} t_j + c_6 \sum_{j=2}^{n} (t_j - 1) + c_7 \sum_{j=2}^{n} (t_j - 1) + c_8(n-1)$$

$$\sum_{j=2}^{n} j = (\sum_{j=1}^{n} j) - 1 = \frac{n(n+1)}{2} - 1$$

Worst Case: List is reversed order

$$\begin{split} T(n) &= c_1 n + c_2 (n-1) + c_4 (n-1) + c_5 \left( \frac{n(n+1)}{2} - 1 \right) + c_6 \left( \frac{(n-1)n}{2} \right) + c_7 \left( \frac{(n-1)n}{2} \right) + c_8 (n-1) \\ &= \left( \frac{c_5}{2} + \frac{c_6}{2} + \frac{c_7}{2} \right) n^2 + \left( c_1 + c_2 + c_4 + \frac{c_5}{2} - \frac{c_6}{2} - \frac{c_7}{2} + c_8 \right) n - (c_2 + c_4 + c_5 + c_8) \\ &= an^2 + bn + c \end{split}$$



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### Selection sort

#### Algorithm 4 Selection Sort

```
1: for i = 1 to n - 1 do
2: min = i
3: for j = i + 1 to n do
4: // Find the index of the i^{th} smallest element
5: if A[j] < A[min] then
6: min = j
7: end if
8: end for
9: Swap A[min] and A[i]
10: end for
```

This yields a running time of

$$\sum_{i=1}^{n-1} n - i = n(n-1) - \sum_{i=1}^{n-1} i = n^2 - n - \frac{n^2 - n}{2} = \frac{n^2 - n}{2} = \Theta(n^2).$$



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### Linear Search

• Time complexity: O(n)

#### Algorithm 4 Linear Search

```
1: procedure Linear(A, n, item)
2: for i \leftarrow 0, n-1 do
3: if A[i] == item then
4: return i
5: end if
6: end for
7: return -1
8: end procedure
```



### References



Introduction to Algorithms, Third Edition, Thomas H. Cormen, Charle E. Leiserson, Ronald L. Rivest, Clifford Stein (clrs).

