

$L^3$

Team Notebook

25/09/2019

CONTENTS

1. Code Templates	1
2. Data Structures	1
2.1. Fenwick Tree	1
2.2. Mergesort Tree	1
2.3. Segment Tree	1
2.4. Sparse Table	2
2.5. Sqrt Decomposition	2
2.6. Treap	2
2.7. Ordered Statistics Tree	3
2.8. Union Find	3
3. Graphs	3
3.1. Single-Source Shortest Paths	3
3.2. All-Pairs Shortest Paths	3
3.3. Strongly Connected Components	3
3.4. Cut Points and Bridges	3
3.5. Biconnected Components	3
3.6. Minimum Spanning Tree	3
3.7. Topological Sorting	3
3.8. Euler Path	3
3.9. Bipartite Matching	3
3.10. Maximum Flow	3
3.11. All-pairs Maximum Flow	4
3.12. Heavy Light Decomposition	4
3.13. Centroid Decomposition	4
3.14. Least Common Ancestor	4
4. Strings	5
4.1. Z-algorithm	5
4.2. Trie	5
4.3. Hashing	5
5. Dynamic Programming	5
5.1. Longest Common Subsequence	5
5.2. Longest Increasing Subsequence	5
5.3. Traveling Salesman	5
6. Mathematics	5
6.1. Special Data Types	5
6.2. Binomial Coefficients	5
6.3. Euclidean Algorithm	5
6.4. Primality Test	5
6.5. Sieve	5
6.6. Phi Function	5
6.7. Modular Exponentiation	5
6.8. Modular Multiplicative Inverse	5
6.9. Chinese Remainder Theorem	5
6.10. Numeric Integration (Simpson's Rule)	5
6.11. Fast Fourier Transform	5
6.12. Josephus Problem	5
6.13. Number of Integer Points Below a Line	5
7. Geometry	5
7.1. Primitives	5

7.2. Lines	5
7.3. Circles	5
7.4. Polygons	5
7.5. Convex Hull (Graham's Scan)	5
7.6. Closest Pair of Points	5
7.7. Rectilinear Minimum Spanning Tree	5
8. Other Algorithms	5
8.1. Coordinate Compression	5
8.2. 2SAT	5
8.3. Nth Permutation	5
8.4. Floyd's Cycle-Finding	5
8.5. Simulated Annealing	5
8.6. Hexagonal Grid Algorithms	5
9. Useful Information (CLEAN THIS UP!!)	5
10. Misc	5
10.1. Debugging Tips	5
10.2. Solution Ideas	5
11. Formulas	5
11.1. Physics	5
11.2. Markov Chains	5
11.3. Burnside's Lemma	5
11.4. Bézout's identity	5
11.5. Misc	5
Practice Contest Checklist	5

1. CODE TEMPLATES

```
#include <bits/stdc++.h> -----//001
typedef long long ll; -----//002
typedef unsigned long long ull; -----//003
typedef std::pair<int, int> ii; -----//004
typedef std::vector<int> vi; -----//005
typedef std::vector<vi> vvi; -----//006
typedef std::vector<ii> vii; -----//007
const int INF = ~(1<<31); -----//008
const ll LINF = (1LL << 60); -----//009
const double EPS = 1e-9; -----//00a
const double pi = acos(-1); -----//00b
```

2. DATA STRUCTURES

2.1. Fenwick Tree.

2.1.1. Point Queries.

```
struct fenwick { -----//00c
- vi ar; -----//00d
- fenwick(vi &ar) : ar(ar.size(), 0) { -----//00e
--- for (int i = 0; i < ar.size(); ++i) { -----//00f
----- ar[i] += _ar[i]; -----//010
----- int j = i | (i+1); -----//011
----- if (j < ar.size()) -----//012
----- ar[j] += ar[i]; } } -----//013
- int sum(int i) { -----//014
--- int res = 0; -----//015
--- for (; i >= 0; i = (i & (i+1)) - 1) -----//016
----- res += ar[i]; -----//017
```

```
--- return res; } -----//018
- int sum(int i, int j) { return sum(j) - sum(i-1); } ----//019
- void add(int i, int val) { -----//01a
--- for (; i < ar.size(); i |= i+1) -----//01b
----- ar[i] += val; } -----//01c
- int get(int i) { -----//01d
--- int res = ar[i]; -----//01e
--- if (i) { -----//01f
----- int lca = (i & (i+1)) - 1; -----//020
----- for (--i; i != lca; i = (i&(i+1))-1) -----//021
----- res -= ar[i]; } -----//022
--- return res; } -----//023
- void set(int i, int val) { add(i, -get(i) + val); } ----//024
- // range update, point query // -----//025
- void add(int i, int j, int val) { -----//026
--- add(i, val); -----//027
--- add(j+1, -val); } -----//028
- int get1(int i) { return sum(i); } -----//029
- ////////////////////////////////////// -----//02a
};; -----//02b
```

2.1.2. Range Queries.

2.2. Mergesort Tree.

2.3. Segment Tree.

2.3.1. Recursive Segment Tree (Point-update).

```
struct segtree { -----//090
- int i, j, val; -----//091
- segtree *l, *r; -----//092
- segtree(int *ar, int _i, int _j) : i(_i), j(_j) { ----//093
--- if (i == j) { -----//094
----- val = ar[i]; -----//095
----- l = r = NULL; -----//096
--- } else { -----//097
----- int k = (i+j) >> 1; -----//098
----- l = new segtree(ar, i, k); -----//099
----- r = new segtree(ar, k+1, j); -----//09a
----- val = l->val + r->val; } } -----//09b
- void update(int _i, int _val) { -----//09c
--- if (i == _i and _i == j) { -----//09d
----- val = _val; -----//09e
--- } else if (_i < i or j < _i) { -----//09f
----- // do nothing -----//0a0
--- } else { -----//0a1
----- l->update(_i, _val); -----//0a2
----- r->update(_i, _val); -----//0a3
----- val = l->val + r->val; } } -----//0a4
- int query(int _i, int _j) { -----//0a5
--- if (_i <= i and j <= _j) { -----//0a6
----- return val; -----//0a7
--- } else if (_j < i or j < _i) { -----//0a8
----- return 0; -----//0a9
--- } else { -----//0aa
----- return l->query(_i, _j) + r->query(_i, _j); } } };
```

2.3.2. *Iterative Segment Tree (Point-update and operation can be non-commutative).*

```

struct segtree { -----//02c
- int n; -----//02d
- int *vals; -----//02e
- segtree(int *ar, int n) { -----//02f
-   this->n = n; -----//030
-   vals = new int[2*n]; -----//031
-   for (int i = 0; i < n; ++i) -----//032
-       vals[n+i] = ar[i]; -----//033
-   for (int i = n-1; i > 0; --i) -----//034
-       vals[i] = vals[i<<1] + vals[i<<1|1]; } -----//035
- void update(int i, int v) { -----//036
-   for (vals[i += n] = v; i > 1; i >=> 1) -----//037
-       vals[i>>1] = vals[i] + vals[i^1]; } -----//038
- int query(int l, int r) { -----//039
-   r++; // without this, the range is [l,r) -----//03a
-   int res = 0; -----//03b
-   for (l += n, r += n; l < r; l >=> 1, r >=> 1) { -----//03c
-       if (l&1) res += vals[l++]; -----//03d
-       if (r&1) res += vals[--r]; } -----//03e
-   return res; } }; -----//03f

```

2.3.3. *Lazy Segment Tree (Range-update).*

```

struct segtree { -----//040
- int i, j, val, temp_val = 0; -----//041
- segtree *l, *r; -----//042
- segtree(int *ar, int _i, int _j) : i(_i), j(_j) { -----//043
-   if (i == j) { -----//044
-       val = ar[i]; -----//045
-       l = r = NULL; -----//046
-   } else { -----//047
-       int k = (i + j) >> 1; -----//048
-       l = new segtree(ar, i, k); -----//049
-       r = new segtree(ar, k+1, j); -----//04a
-       val = l->val + r->val; } } -----//04b
- void visit() { -----//04c
-   if (temp_val) { -----//04d
-       val += (j-i+1) * temp_val; -----//04e
-       if (l) { -----//04f
-           l->temp_val += temp_val; -----//050
-           r->temp_val += temp_val; } -----//051
-       temp_val = 0; } } -----//052
- void increase(int _i, int _j, int _inc) { -----//053
-   visit(); -----//054
-   if (_i <= i && j <= _j) { -----//055
-       temp_val += _inc; -----//056
-       visit(); -----//057
-   } else if (_j < i or j < _i) { -----//058
-       // do nothing -----//059
-   } else { -----//05a
-       l->increase(_i, _j, _inc); -----//05b
-       r->increase(_i, _j, _inc); -----//05c
-       val = l->val + r->val; } } -----//05d
- int query(int _i, int _j) { -----//05e
-   visit(); -----//05f

```

```

-   if (_i <= i and j <= _j) { -----//060
-       return val; -----//061
-   } else if (_j < i || j < _i) { -----//062
-       return 0; -----//063
-   } else { -----//064
-       return l->query(_i, _j) + r->query(_i, _j); } } };

```

2.3.4. *Persistent Segmentr Tree (Point-update).*

```

struct node { int l, r, lid, rid, val; }; -----//066
struct segtree { -----//067
-   node *nodes; -----//068
-   int n, node_cnt = 0; -----//069
-   segtree(int n, int capacity) { -----//06a
-       this->n = n; -----//06b
-       nodes = new node[capacity]; } -----//06c
-   int build (int *ar, int l, int r) { -----//06d
-       if (l > r) return -1; -----//06e
-       int id = node_cnt++; -----//06f
-       nodes[id].l = l; -----//070
-       nodes[id].r = r; -----//071
-       if (l == r) { -----//072
-           nodes[id].lid = -1; -----//073
-           nodes[id].rid = -1; -----//074
-           nodes[id].val = ar[l]; -----//075
-       } else { -----//076
-           int m = (l + r) / 2; -----//077
-           nodes[id].lid = build(ar, l, m); -----//078
-           nodes[id].rid = build(ar, m+1, r); -----//079
-           nodes[id].val = nodes[nodes[id].lid].val + -----//07a
-               nodes[nodes[id].rid].val; } -----//07b
-       return id; } -----//07c
-   int update(int id, int idx, int delta) { -----//07d
-       if (id == -1) -----//07e
-           return -1; -----//07f
-       if (idx < nodes[id].l or nodes[id].r < idx) -----//080
-           return id; -----//081
-       int nid = node_cnt++; -----//082
-       nodes[nid].l = nodes[id].l; -----//083
-       nodes[nid].r = nodes[id].r; -----//084
-       nodes[nid].lid = update(nodes[id].lid, idx, delta); -----//085
-       nodes[nid].rid = update(nodes[id].rid, idx, delta); -----//086
-       nodes[nid].val = nodes[id].val + delta; -----//087
-       return nid; } -----//088
-   int query(int id, int l, int r) { -----//089
-       if (r < nodes[id].l or nodes[id].r < l) -----//08a
-           return 0; -----//08b
-       if (l <= nodes[id].l and nodes[id].r <= r) -----//08c
-           return nodes[id].val; -----//08d
-       return query(nodes[id].lid, l, r) + -----//08e
-           query(nodes[id].rid, l, r); } }; -----//08f

```

2.4. *Sparse Table.*

2.5. *Sqrt Decomposition.*

2.6. *Treap.*

2.6.1. *Explicit Treap.*

2.6.2. *Implicit Treap.*

```

struct cartree { -----//0ac
-   typedef struct _Node { -----//0ad
-       int node_val, subtree_val, delta, prio, size; -----//0ae
-       _Node *l, *r; -----//0af
-       _Node(int val) : node_val(val), subtree_val(val), -----//0b0
-           delta(0), prio((rand()<<16)^rand()), size(1), -----//0b1
-           l(NULL), r(NULL) {} -----//0b2
-   } _Node(); delete l; delete r; } -----//0b3
-   } *Node; -----//0b4
-   int get_subtree_val(Node v) { -----//0b5
-       return v ? v->subtree_val : 0; } -----//0b6
-   int get_size(Node v) { return v ? v->size : 0; } -----//0b7
-   void apply_delta(Node v, int delta) { -----//0b8
-       if (!v) return; -----//0b9
-       v->delta += delta; -----//0ba
-       v->node_val += delta; -----//0bb
-       v->subtree_val += delta * get_size(v); } -----//0bc
-   void push_delta(Node v) { -----//0bd
-       if (!v) return; -----//0be
-       apply_delta(v->l, v->delta); -----//0bf
-       apply_delta(v->r, v->delta); -----//0c0
-       v->delta = 0; } -----//0c1
-   void update(Node v) { -----//0c2
-       if (!v) return; -----//0c3
-       v->subtree_val = get_subtree_val(v->l) + v->node_val -----//0c4
-           + get_subtree_val(v->r); -----//0c5
-       v->size = get_size(v->l) + 1 + get_size(v->r); } -----//0c6
-   Node merge(Node l, Node r) { -----//0c7
-       push_delta(l); push_delta(r); -----//0c8
-       if (!l || !r) return l ? l : r; -----//0c9
-       if (l->size <= r->size) { -----//0ca
-           l->r = merge(l->r, r); -----//0cb
-           update(l); -----//0cc
-           return l; -----//0cd
-       } else { -----//0ce
-           r->l = merge(l, r->l); -----//0cf
-           update(r); -----//0d0
-           return r; } } -----//0d1
-   void split(Node v, int key, Node &l, Node &r) { -----//0d2
-       push_delta(v); -----//0d3
-       l = r = NULL; -----//0d4
-       if (!v) return; -----//0d5
-       if (key <= get_size(v->l)) { -----//0d6
-           split(v->l, key, l, v->l); -----//0d7
-           r = v; -----//0d8
-       } else { -----//0d9
-           split(v->r, key - get_size(v->l) - 1, v->r, r); -----//0da
-           l = v; } -----//0db
-       update(v); } -----//0dc
-   Node root; -----//0dd
public: -----//0de
-   cartree() : root(NULL) {} -----//0df
-   ~cartree() { delete root; } -----//0e0
-   int get(Node v, int key) { -----//0e1
-       push_delta(v); -----//0e2

```

```
--- if (key < get_size(v->l)) -----//0e3
--- return get(v->l, key); -----//0e4
--- else if (key > get_size(v->l)) -----//0e5
--- return get(v->r, key - get_size(v->l) - 1); -----//0e6
--- return v->node_val; } -----//0e7
- int get(int key) { return get(root, key); } -----//0e8
- void insert(Node item, int key) { -----//0e9
- Node l, r; -----//0ea
- split(root, key, l, r); -----//0eb
- root = merge(merge(l, item), r); } -----//0ec
- void insert(int key, int val) { -----//0ed
- insert(new _Node(val), key); } -----//0ee
- void erase(int key) { -----//0ef
- Node l, m, r; -----//0f0
- split(root, key + 1, m, r); -----//0f1
- split(m, key, l, m); -----//0f2
- delete m; -----//0f3
- root = merge(l, r); } -----//0f4
- int query(int a, int b) { -----//0f5
- Node l1, r1; -----//0f6
- split(root, b+1, l1, r1); -----//0f7
- Node l2, r2; -----//0f8
- split(l1, a, l2, r2); -----//0f9
- int res = get_subtree_val(r2); -----//0fa
- l1 = merge(l2, r2); -----//0fb
- root = merge(l1, r1); -----//0fc
- return res; } -----//0fd
- int update(int a, int b, int delta) { -----//0fe
- Node l1, r1; -----//0ff
- split(root, b+1, l1, r1); -----//100
- Node l2, r2; -----//101
- split(l1, a, l2, r2); -----//102
- apply_delta(r2, delta); -----//103
- l1 = merge(l2, r2); -----//104
- root = merge(l1, r1); } -----//105
- int size() { return get_size(root); } }; -----//106
```

2.6.3. Persistent Treap.

2.7. Ordered Statistics Tree.

2.8. Union Find.

```
struct union_find { -----//107
- vi p; union_find(int n) : p(n, -1) { } -----//108
- int find(int x) { return p[x] < 0 ? x : p[x] = find(p[x]); }
- bool unite(int x, int y) { -----//10a
- int xp = find(x), yp = find(y); -----//10b
- if (xp == yp) return false; -----//10c
- if (p[xp] > p[yp]) swap(xp,yp); -----//10d
- p[xp] += p[yp], p[yp] = xp; -----//10e
- return true; } -----//10f
- int size(int x) { return -p[find(x)]; } }; -----//110
```

3. GRAPHS

Using adjacency list:

```
struct graph { -----//1a7
- int n; -----//1a8
```

```
    vii *adj; -----//1a9
    int *dist; -----//1aa
    graph(int n) { -----//1ab
    this->n = n; -----//1ac
    adj = new vii[n]; -----//1ad
    dist = new int[n]; } -----//1ae
    void add_edge(int u, int v, int w) { -----//1af
    adj[u].push_back({v, w}); -----//1b0
    /*adj[v].push_back({u, w});*/ } }; -----//1b1

    Using adjacency matrix:

    struct graph { -----//1b2
    int n; -----//1b3
    int **mat; -----//1b4
    graph(int n) { -----//1b5
    this->n = n; -----//1b6
    mat = new int*[n]; -----//1b7
    for (int i = 0; i < n; ++i) { -----//1b8
    mat[i] = new int[n]; -----//1b9
    for (int j = 0; j < n; ++j) -----//1ba
    mat[i][j] = INF; -----//1bb
    mat[i][i] = 0; } } -----//1bc
    void add_edge(int u, int v, int w) { -----//1bd
    mat[u][v] = std::min(mat[u][v], w); -----//1be
    /*mat[v][u] = std::min(mat[v][u], w);*/ } }; -----//1bf

    Using edge list:
```

```
struct edge { -----//1c0
    int u, v, w; -----//1c1
    edge(int u, int v, int w) : u(u), v(v), w(w) {} -----//1c2
    const bool operator <(const edge &other) const { -----//1c3
    return w < other.w; } }; -----//1c4
    struct graph { -----//1c5
    int n; -----//1c6
    std::vector<edge> edges; -----//1c7
    graph(int n) : n(n) {} -----//1c8
    void add_edge(int u, int v, int w) { -----//1c9
    edges.push_back(edge(u, v, w)); } }; -----//1ca
```

3.1. Single-Source Shortest Paths.

3.1.1. Dijkstra.

```
#include "graph_template_adjlist.cpp" -----//122
void dijkstra(int s, int n, int *dist, vii *adj) { -----//123
- for (int u = 0; u < n; ++u) -----//124
- dist[u] = INF; -----//125
- dist[s] = 0; -----//126
- std::priority_queue<ii, vii, std::greater<ii> > pq; -----//127
- pq.push({0, s}); -----//128
- while (!pq.empty()) { -----//129
- int u = pq.top().second; -----//12a
- int d = pq.top().first; -----//12b
- pq.pop(); -----//12c
- if (dist[u] < d) -----//12d
- continue; -----//12e
- dist[u] = d; -----//12f
- for (auto &e : adj[u]) { -----//130
- int v = e.first; -----//131
```

```
----- int w = e.second; -----//132
----- if (dist[v] > dist[u] + w) { -----//133
----- dist[v] = dist[u] + w; -----//134
----- pq.push({dist[v], v}); } } } } -----//135
```

3.1.2. Bellman-Ford.

```
#include "graph_template_adjlist.cpp" -----//111
void bellman_ford(int s, int n, int *dist, vii *adj) { ---//112
- for (int u = 0; u < n; ++u) -----//113
- dist[u] = INF; -----//114
- dist[s] = 0; -----//115
- for (int i = 0; i < n-1; ++i) -----//116
- for (int u = 0; u < n; ++u) -----//117
- for (auto &e : adj[u]) -----//118
- if (dist[u] + e.second < dist[e.first]) -----//119
- dist[e.first] = dist[u] + e.second; } -----//11a
// you can call this after running bellman_ford() -----//11b
bool has_neg_cycle(int n, int *dist, vii *adj) { -----//11c
- for (int u = 0; u < n; ++u) -----//11d
- for (auto &e : adj[u]) -----//11e
- if (dist[e.first] > dist[u] + e.second) -----//11f
- return true; -----//120
- return false; } -----//121
```

3.2. All-Pairs Shortest Paths.

3.2.1. Floyd-Washall.

```
#include "graph_template_adjmat.cpp" -----//1a0
void floyd_warshall(int n, int **mat) { -----//1a1
- for (int k = 0; k < n; ++k) -----//1a2
- for (int i = 0; i < n; ++i) -----//1a3
- for (int j = 0; j < n; ++j) -----//1a4
- if (mat[i][k] + mat[k][j] < mat[i][j]) -----//1a5
- mat[i][j] = mat[i][k] + mat[k][j]; } -----//1a6
```

3.3. Strongly Connected Components.

3.3.1. Kosaraju.

3.4. Cut Points and Bridges.

3.5. Biconnected Components.

3.5.1. Bridge Tree.

3.5.2. Block-Cut Tree.

3.6. Minimum Spanning Tree.

3.6.1. Kruskal.

3.6.2. Prim.

3.7. Topological Sorting.

3.8. Euler Path.

3.9. Bipartite Matching.

3.9.1. Alternating Paths Algorithm.

3.9.2. Hopcroft-Karp Algorithm.

3.10. Maximum Flow.

3.10.1. Edmonds-Karp.

```
struct max_flow { -----//173
- int n, s, t, *par, **c, **f; -----//174
- vi *adj; -----//175
- max_flow(int n, int s, int t) : n(n), s(s), t(t) { -----//176
-- adj = new std::vector<int>[n]; -----//177
-- par = new int[n]; -----//178
-- c = new int*[n]; -----//179
-- f = new int*[n]; -----//17a
-- for (int i = 0; i < n; ++i) { -----//17b
--     c[i] = new int[n]; -----//17c
--     f[i] = new int[n]; -----//17d
--     for (int j = 0; j < n; ++j) -----//17e
--         c[i][j] = f[i][j] = 0; } } -----//17f
- void add_edge(int u, int v, int w) { -----//180
-- adj[u].push_back(v); -----//181
-- adj[v].push_back(u); -----//182
-- c[u][v] += w; } -----//183
- int res(int i, int j) { return c[i][j] - f[i][j]; } -----//184
- bool bfs() { -----//185
-- std::queue<int> q; -----//186
-- q.push(this->s); -----//187
-- while (!q.empty()) { -----//188
--     int u = q.front(); q.pop(); -----//189
--     for (int v : adj[u]) { -----//18a
--         if (res(u, v) > 0 and par[v] == -1) { -----//18b
--             par[v] = u; -----//18c
--             if (v == this->t) -----//18d
--                 return true; -----//18e
--             q.push(v); } } -----//18f
-- return false; } -----//190
- bool aug_path() { -----//191
-- for (int u = 0; u < n; ++u) -----//192
--     par[u] = -1; -----//193
-- par[s] = s; -----//194
-- return bfs(); } -----//195
- int calc_max_flow() { -----//196
-- int ans = 0; -----//197
-- while (aug_path()) { -----//198
--     int flow = INF; -----//199
--     for (int u = t; u != s; u = par[u]) -----//19a
--         flow = std::min(flow, res(par[u], u)); -----//19b
--     for (int u = t; u != s; u = par[u]) -----//19c
--         f[par[u]][u] += flow, f[u][par[u]] -= flow; -----//19d
--     ans += flow; } -----//19e
-- return ans; } }; -----//19f
```

3.10.2. Dinic.

```
struct max_flow { -----//136
- int n, s, t, *adj_ptr, *dist, *par, **c, **f; -----//137
- vi *adj; -----//138
- max_flow(int n, int s, int t) : n(n), s(s), t(t) { -----//139
-- adj = new std::vector<int>[n]; -----//13a
-- adj_ptr = new int[n]; -----//13b
-- dist = new int[n]; -----//13c
-- par = new int[n]; -----//13d
```

```
-- c = new int*[n]; -----//13e
-- f = new int*[n]; -----//13f
-- for (int i = 0; i < n; ++i) { -----//140
--     c[i] = new int[n]; -----//141
--     f[i] = new int[n]; -----//142
--     for (int j = 0; j < n; ++j) -----//143
--         c[i][j] = f[i][j] = 0; } } -----//144
- void add_edge(int u, int v, int w) { -----//145
-- adj[u].push_back(v); -----//146
-- adj[v].push_back(u); -----//147
-- c[u][v] += w; } -----//148
- int res(int i, int j) { return c[i][j] - f[i][j]; } -----//149
- void reset(int *ar, int val) { -----//14a
-- for (int i = 0; i < n; ++i) -----//14b
--     ar[i] = val; } -----//14c
- bool make_level_graph() { -----//14d
-- reset(dist, -1); -----//14e
-- std::queue<int> q; -----//14f
-- q.push(s); -----//150
-- dist[s] = 0; -----//151
-- while (!q.empty()) { -----//152
--     int u = q.front(); q.pop(); -----//153
--     for (int v : adj[u]) { -----//154
--         if (res(u, v) > 0 and dist[v] == -1) { -----//155
--             dist[v] = dist[u] + 1; -----//156
--             q.push(v); } } } -----//157
-- return dist[t] != -1; } -----//158
- bool next(int u, int v) { return dist[v] == dist[u] + 1; } -----//15a
- bool dfs(int u) { -----//15b
-- if (u == t) return true; -----//15c
-- for (int &i = adj_ptr[u]; i < adj[u].size(); ++i) { -----//15d
--     int v = adj[u][i]; -----//15e
--     if (next(u, v) and res(u, v) > 0 and dfs(v)) { -----//15f
--         par[v] = u; -----//160
--         return true; } } -----//161
-- dist[u] = -1; -----//162
-- return false; } -----//163
- bool aug_path() { -----//164
-- reset(par, -1); -----//165
-- par[s] = s; -----//166
-- return dfs(s); } -----//167
- int calc_max_flow() { -----//168
-- int ans = 0; -----//169
-- while (make_level_graph()) { -----//16a
--     reset(adj_ptr, 0); -----//16b
--     while (aug_path()) { -----//16c
--         int flow = INF; -----//16d
--         for (int u = t; u != s; u = par[u]) -----//16e
--             flow = std::min(flow, res(par[u], u)); -----//16f
--         for (int u = t; u != s; u = par[u]) -----//170
--             f[par[u]][u] += flow, f[u][par[u]] -= flow; -----//171
--         ans += flow; } } -----//172
-- return ans; } };
```

3.11. All-pairs Maximum Flow.

3.11.1. Gomory-Hu.

3.12. Heavy Light Decomposition.

```
#include "segment_tree.cpp" -----//1cb
struct heavy_light_tree { -----//1cc
- int n; -----//1cd
- std::vector<int> *adj; -----//1ce
- segtree *segment_tree; -----//1cf
- int *par, *heavy, *dep, *path_root, *pos; -----//1d0
- heavy_light_tree(int n) { -----//1d1
-- this->n = n; -----//1d2
-- this->adj = new std::vector<int>[n]; -----//1d3
-- segment_tree = new segtree(0, n-1); -----//1d4
-- par = new int[n]; -----//1d5
-- heavy = new int[n]; -----//1d6
-- dep = new int[n]; -----//1d7
-- path_root = new int[n]; -----//1d8
-- pos = new int[n]; } -----//1d9
- void add_edge(int u, int v) { -----//1da
-- adj[u].push_back(v); -----//1db
-- adj[v].push_back(u); } -----//1dc
- void build(int root) { -----//1dd
-- for (int u = 0; u < n; ++u) -----//1de
--     heavy[u] = -1; -----//1df
-- par[root] = root; -----//1e0
-- dep[root] = 0; -----//1e1
-- dfs(root); -----//1e2
-- for (int u = 0, p = 0; u < n; ++u) { -----//1e3
--     if (par[u] == -1 or heavy[par[u]] != u) { -----//1e4
--         for (int v = u; v != -1; v = heavy[v]) { -----//1e5
--             path_root[v] = u; -----//1e6
--             pos[v] = p++; } } } -----//1e7
- int dfs(int u) { -----//1e8
-- int sz = 1; -----//1e9
-- int max_subtree_sz = 0; -----//1ea
-- for (int v : adj[u]) { -----//1eb
--     if (v != par[u]) { -----//1ec
--         par[v] = u; -----//1ed
--         dep[v] = dep[u] + 1; -----//1ee
--         int subtree_sz = dfs(v); -----//1ef
--         if (max_subtree_sz < subtree_sz) { -----//1f0
--             max_subtree_sz = subtree_sz; -----//1f1
--             heavy[u] = v; } -----//1f2
--         sz += subtree_sz; } } -----//1f3
-- return sz; } -----//1f4
- int query(int u) { -----//1f5
-- return segment_tree->sum(pos[u], pos[u]); } -----//1f6
- void update(int u, int v, int c) { -----//1f7
-- for (; path_root[u] != path_root[v]; -----//1f8
--     v = par[path_root[v]]) { -----//1f9
--     if (dep[path_root[u]] > dep[path_root[v]]) -----//1fa
--         std::swap(u, v); -----//1fb
--     segment_tree->increase(pos[path_root[v]], pos[v], c); }
-- segment_tree->increase(pos[u], pos[v], c); } };
```

3.13. Centroid Decomposition.

3.14. Least Common Ancestor.

3.14.1. *Binary Lifting.*

3.14.2. *Tarjan’s Offline Algorithm.*

4. STRINGS

4.1. **Z-algorithm.**

4.2. **Trie.**

4.3. **Hashing.**

5. DYNAMIC PROGRAMMING

5.1. **Longest Common Subsequence.**

5.2. **Longest Increasing Subsequence.**

5.3. **Traveling Salesman.**

6. MATHEMATICS

6.1. **Special Data Types.**

6.1.1. *Fraction.*

6.1.2. *BigInteger.*

6.1.3. *Matrix.*

6.1.4. *Dates.*

6.2. **Binomial Coefficients.**

6.3. **Euclidean Algorithm.**

6.4. **Primality Test.**

6.4.1. *Optimized Brute Force.*

6.4.2. *Miller-Rabin.*

6.4.3. *Pollard’s Rho Algorithm.*

6.5. **Sieve.**

6.5.1. *Sieve of Eratosthenes.*

6.5.2. *Divisor Sieve (Modified Sieve of Eratosthenes).*

6.5.3. *Phi Sieve.*

6.6. **Phi Function.**

6.7. **Modular Exponentiation.**

6.8. **Modular Multiplicative Inverse.**

6.9. **Chinese Remainder Theorem.**

6.10. **Numeric Integration (Simpson’s Rule).**

6.11. **Fast Fourier Transform.**

6.12. **Josephus Problem.**

6.13. **Number of Integer Points Below a Line.**

7. GEOMETRY

7.1. **Primitives.**

7.2. **Lines.**

7.3. **Circles.**

7.4. **Polygons.**

7.5. **Convex Hull (Graham’s Scan).**

7.6. **Closest Pair of Points.**

7.7. **Rectilinear Minimum Spanning Tree.**

8. OTHER ALGORITHMS

8.1. **Coordinate Compression.**

8.2. **2SAT.**

8.3. **Nth Permutation.**

8.4. **Floyd’s Cycle-Finding.**

8.5. **Simulated Annealing.**

8.6. **Hexagonal Grid Algorithms.**

· sufficient: QI and  $C[b][c] \leq C[a][d]$ ,  $a \leq b \leq c \leq d$

- Permutations
  - \* Consider the cycles of the permutation
- Functions
  - \* Sum of piecewise-linear functions is a piecewise-linear function
  - \* Sum of convex (concave) functions is convex (concave)
- Modular arithmetic
  - \* Chinese Remainder Theorem
  - \* Linear Congruence
- Sieve
- System of linear equations
- Values too big to represent?
  - \* Compute using the logarithm
  - \* Divide everything by some large value
- Linear programming
  - \* Is the dual problem easier to solve?
- Can the problem be modeled as a different combinatorial problem? Does that simplify calculations?
- Logic
  - 2-SAT
  - XOR-SAT (Gauss elimination or Bipartite matching)
- Meet in the middle
- Only work with the smaller half ( $\log(n)$ )
- Strings
  - Trie (maybe over something weird, like bits)
  - Suffix array
  - Suffix automaton (+DP?)
  - Aho-Corasick
  - Suffix Tree
- Work with  $S + S$
- Cashing
- Traveling salesman tour, tree to array
- Segment trees
  - Lazy propagation
  - Persistent
  - Implicit
  - Segment tree of X
- Computational geometry
  - Minkowski sum (of convex sets)
  - Rotating calipers
  - Sweep line (horizontally or vertically?)
  - Sweep angle
  - Convex hull
- Fix a parameter (possibly the answer).
- Are there few distinct values?
- Binary search
- Sliding Window (+ Monotonic Queue)
- Computing a Convolution? Fast Fourier Transform
- Computing a 2D Convolution? FFT on each row, and then on each column
- Exact Cover (+ Algorithm X)
- Cycle-Finding
  - What is the smallest set of values that identify the solution? The cycle structure of the permutation? The powers of primes in the factorization?
  - Look at the complement problem

- Stack overflow? Recursive DFS on tree that is actually a long path?
- Floating-point numbers
  - Getting **NaN**? Make sure **acos** etc. are not getting values out of their range (perhaps **1+eps**).
  - Rounding negative numbers?
  - Outputting in scientific notation?
- Wrong Answer?
  - Read the problem statement again!
  - Are multiple test cases being handled correctly? Try repeating the same test case many times.
  - Integer overflow?
  - Think very carefully about boundaries of all input parameters
  - Try out possible edge cases:
    - \*  $n = 0, n = -1, n = 1, n = 2^{31} - 1$  or  $n = -2^{31}$
    - \* List is empty, or contains a single element
    - \*  $n$  is even,  $n$  is odd
    - \* Graph is empty, or contains a single vertex
    - \* Graph is a multigraph (loops or multiple edges)
    - \* Polygon is concave or non-simple
  - Is initial condition wrong for small cases?
  - Are you sure the algorithm is correct?
  - Explain your solution to someone.
  - Are you using any functions that you don't completely understand? Maybe STL functions?
  - Maybe you (or someone else) should rewrite the solution?
  - Can the input line be empty?
- Run-Time Error?
  - Is it actually Memory Limit Exceeded?

- Greedy
- Randomized
- Optimizations
  - Use bitset (/64)
  - Switch order of loops (cache locality)
- Process queries offline
  - Mo's algorithm
- Square-root decomposition
- Precomputation
- Efficient simulation
  - Mo's algorithm
  - Sqrt decomposition
  - Store  $2^k$  jump pointers
- Data structure techniques
  - Sqrt buckets
  - Store  $2^k$  jump pointers
  - $2^k$  merging trick
- Counting
  - Inclusion-exclusion principle
  - Generating functions
- Graphs
  - Can we model the problem as a graph?
  - Can we use any properties of the graph?
  - Strongly connected components
  - Cycles (or odd cycles)
  - Bipartite (no odd cycles)
    - \* Bipartite matching
    - \* Hall's marriage theorem
    - \* Stable Marriage
  - Cut vertex/bridge
  - Biconnected components
  - Degrees of vertices (odd/even)
  - Trees
    - \* Heavy-light decomposition
    - \* Centroid decomposition
    - \* Least common ancestor
    - \* Centers of the tree
  - Eulerian path/circuit
  - Chinese postman problem
  - Topological sort
  - (Min-Cost) Max Flow
  - Min Cut
    - \* Maximum Density Subgraph
  - Huffman Coding
  - Min-Cost Arborescence
  - Steiner Tree
  - Kirchoff's matrix tree theorem
  - Prüfer sequences
  - Lovász Toggle
  - Look at the DFS tree (which has no cross-edges)
  - Is the graph a DFA or NFA?
    - \* Is it the Synchronizing word problem?
- Mathematics
  - Is the function multiplicative?
  - Look for a pattern

- Minimize something instead of maximizing
- Immediately enforce necessary conditions. (All values greater than 0? Initialize them all to 1)
- Add large constant to negative numbers to make them positive
- Counting/ Bucket sort

11. FORMULAS

- **Legendre symbol:**  $(\frac{a}{b}) = a^{(b-1)/2} \pmod{b}$ ,  $b$  odd prime.
- **Heron’s formula:** A triangle with side lengths  $a, b, c$  has area  $\sqrt{s(s-a)(s-b)(s-c)}$  where  $s = \frac{a+b+c}{2}$ .
- **Pick’s theorem:** A polygon on an integer grid strictly containing  $i$  lattice points and having  $b$  lattice points on the boundary has area  $i + \frac{b}{2} - 1$ . (Nothing similar in higher dimensions)
- **Euler’s totient:** The number of integers less than  $n$  that are coprime to  $n$  are  $n \prod_{p|n} \left(1 - \frac{1}{p}\right)$  where each  $p$  is a distinct prime factor of  $n$ .
- **König’s theorem:** In any bipartite graph  $G = (L \cup R, E)$ , the number of edges in a maximum matching is equal to the number of vertices in a minimum vertex cover. Let  $U$  be the set of unmatched vertices in  $L$ , and  $Z$  be the set of vertices that are either in  $U$  or are connected to  $U$  by an alternating path. Then  $K = (L \setminus Z) \cup (R \cap Z)$  is the minimum vertex cover.
- A minium Steiner tree for  $n$  vertices requires at most  $n-2$  additional Steiner vertices.
- The number of vertices of a graph is equal to its minimum vertex cover number plus the size of a maximum independent set.
- **Lagrange polynomial** through points  $(x_0, y_0), \dots, (x_k, y_k)$  is  $L(x) = \sum_{j=0}^k y_j \prod_{\substack{0 \leq m \leq k \\ m \neq j}} \frac{x-x_m}{x_j-x_m}$
- **Hook length formula:** If  $\lambda$  is a Young diagram and  $h_\lambda(i, j)$  is the hook-length of cell  $(i, j)$ , then then the number of Young tableaux  $d_\lambda = n! / \prod h_\lambda(i, j)$ .
- **Möbius inversion formula:** If  $f(n) = \sum_{d|n} g(d)$ , then  $g(n) = \sum_{d|n} \mu(d) f(n/d)$ . If  $f(n) = \sum_{m=1}^n g(\lfloor n/m \rfloor)$ , then  $g(n) = \sum_{m=1}^n \mu(m) f(\lfloor \frac{n}{m} \rfloor)$ .
- #primitive pythagorean triples with hypotenuse  $< n$  approx  $n/(2\pi)$ .
- **Frobenius Number:** largest number which can’t be expressed as a linear combination of numbers  $a_1, \dots, a_n$  with non-negative coefficients.  $g(a_1, a_2) = a_1 a_2 - a_1 - a_2$ ,  $N(a_1, a_2) = (a_1 - 1)(a_2 - 1)/2$ .  $g(d \cdot a_1, d \cdot a_2, a_3) = d \cdot g(a_1, a_2, a_3) + a_3(d - 1)$ . An integer  $x > (\max_i a_i)^2$  can be expressed in such a way iff.  $x \mid \gcd(a_1, \dots, a_n)$ .

11.1. Physics.

- **Snell’s law:**  $\frac{\sin \theta_1}{v_1} = \frac{\sin \theta_2}{v_2}$

11.2. **Markov Chains.** A Markov Chain can be represented as a weighted directed graph of states, where the weight of an edge represents the probability of transitioning over that edge in one timestep. Let  $P^{(m)} = (p_{ij}^{(m)})$  be the probability matrix of transitioning from state  $i$  to state  $j$  in  $m$  timesteps, and note that  $P^{(1)}$  is the adjacency matrix of the graph. **Chapman-Kolmogorov:**  $p_{ij}^{(m+n)} = \sum_k p_{ik}^{(m)} p_{kj}^{(n)}$ . It follows that  $P^{(m+n)} = P^{(m)} P^{(n)}$  and  $P^{(m)} = P^m$ . If  $p^{(0)}$  is the initial probability distribution (a vector), then  $p^{(0)} P^{(m)}$  is the probability distribution after  $m$  timesteps.

The return times of a state  $i$  is  $R_i = \{m \mid p_{ii}^{(m)} > 0\}$ , and  $i$  is *aperiodic* if  $\gcd(R_i) = 1$ . A MC is aperiodic if any of its vertices is aperiodic. A MC is *irreducible* if the corresponding graph is strongly connected.

A distribution  $\pi$  is stationary if  $\pi P = \pi$ . If MC is irreducible then  $\pi_i = 1/\mathbb{E}[T_i]$ , where  $T_i$  is the expected time between two visits at  $i$ .  $\pi_j/\pi_i$  is the expected number of visits at  $j$  in between two consecutive visits at  $i$ . A MC is *ergodic* if  $\lim_{m \rightarrow \infty} p^{(0)} P^m = \pi$ . A MC is ergodic iff. it is irreducible and aperiodic.

A MC for a random walk in an undirected weighted graph (unweighted graph can be made weighted by adding 1-weights) has  $p_{uv} = w_{uv} / \sum_x w_{ux}$ . If the graph is connected, then  $\pi_u = \sum_x w_{ux} / \sum_v \sum_x w_{vx}$ . Such a random walk is aperiodic iff. the graph is not bipartite.

An *absorbing* MC is of the form  $P = \begin{pmatrix} Q & R \\ 0 & I_r \end{pmatrix}$ . Let  $N = \sum_{m=0}^\infty Q^m = (I_t - Q)^{-1}$ . Then, if starting in state  $i$ , the expected number of steps till absorption is the  $i$ -th entry in  $N1$ . If starting in state  $i$ , the probability of being absorbed in state  $j$  is the  $(i, j)$ -th entry of  $NR$ . Many problems on MC can be formulated in terms of a system of recurrence relations, and then solved using Gaussian elimination.

11.3. **Burnside’s Lemma.** Let  $G$  be a finite group that acts on a set  $X$ . For each  $g$  in  $G$  let  $X^g$  denote the set of elements in  $X$  that are fixed by  $g$ . Then the number of orbits

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$$

$$Z(S_n) = \frac{1}{n} \sum_{l=1}^n a_l Z(S_{n-l})$$

11.4. **Bézout’s identity.** If  $(x, y)$  is any solution to  $ax + by = d$  (e.g. found by the Extended Euclidean Algorithm), then all solutions are given by

$$\left(x + k \frac{b}{\gcd(a, b)}, y - k \frac{a}{\gcd(a, b)}\right)$$

11.5. Misc.

11.5.1. *Determinants and PM.*

$$\det(A) = \sum_{\sigma \in S_n} \text{sgn}(\sigma) \prod_{i=1}^n a_{i, \sigma(i)}$$

$$\text{perm}(A) = \sum_{\sigma \in S_n} \prod_{i=1}^n a_{i, \sigma(i)}$$

$$\begin{aligned} pf(A) &= \frac{1}{2^n n!} \sum_{\sigma \in S_{2n}} \text{sgn}(\sigma) \prod_{i=1}^n a_{\sigma(2i-1), \sigma(2i)} \\ &= \sum_{M \in \text{PM}(n)} \text{sgn}(M) \prod_{(i, j) \in M} a_{i, j} \end{aligned}$$

11.5.2. *BEST Theorem.* Count directed Eulerian cycles. Number of OST given by Kirchoff’s Theorem (remove r/c with root)  $\# \text{OST}(G, r) \cdot \prod_v (d_v - 1)!$

11.5.3. *Primitive Roots.* Only exists when  $n$  is  $2, 4, p^k, 2p^k$ , where  $p$  odd prime. Assume  $n$  prime. Number of primitive roots  $\phi(\phi(n))$  Let  $g$  be primitive root. All primitive roots are of the form  $g^k$  where  $k, \phi(p)$  are coprime.

$k$ -roots:  $g^{i \cdot \phi(n)/k}$  for  $0 \leq i < k$

11.5.4. *Sum of primes.* For any multiplicative  $f$ :

$$S(n, p) = S(n, p-1) - f(p) \cdot (S(n/p, p-1) - S(p-1, p-1))$$

11.5.5. *Floor.*

$$\lfloor \lfloor x/y \rfloor / z \rfloor = \lfloor x/(yz) \rfloor$$

$$x \% y = x - y \lfloor x/y \rfloor$$

PRACTICE CONTEST CHECKLIST

- How many operations per second? Compare to local machine.
- What is the stack size?
- How to use printf/scanf with long long/long double?
- Are `__int128` and `__float128` available?
- Does MLE give RTE or MLE as a verdict? What about stack overflow?
- What is `RAND_MAX`?
- How does the judge handle extra spaces (or missing newlines) in the output?
- Look at documentation for programming languages.
- Try different programming languages: C++, Java and Python.
- Try the submit script.
- Try local programs: `i?python[23]`, `factor`.
- Try submitting with `assert(false)` and `assert(true)`.
- Return-value from `main`.
- Look for directory with sample test cases.
- Make sure printing works.
- Remove this page from the notebook.