

AdMU Progar

Team Notebook

06/11/2019

CONTENTS

| | |
|---|----|
| 1. Code Templates | 1 |
| 2. Data Structures | 1 |
| 2.1. Union Find | 1 |
| 2.2. Segment Tree | 1 |
| 2.3. Fenwick Tree | 3 |
| 2.4. Treap | 3 |
| 2.5. Splay Tree | 4 |
| 2.6. Ordered Statistics Tree | 5 |
| 2.7. Sparse Table | 5 |
| 3. Graphs | 5 |
| 3.1. Single-Source Shortest Paths | 5 |
| 3.2. All-Pairs Shortest Paths | 6 |
| 3.3. Strongly Connected Components | 6 |
| 3.4. Minimum Mean Weight Cycle | 6 |
| 3.5. Cut Points and Bridges | 6 |
| 3.6. Biconnected Components | 6 |
| 3.7. Minimum Spanning Tree | 6 |
| 3.8. Euler Path/Cycle | 7 |
| 3.9. Bipartite Matching | 7 |
| 3.10. Maximum Flow | 7 |
| 3.11. All-pairs Maximum Flow | 8 |
| 3.12. Minimum Arborescence | 9 |
| 3.13. Blossom algorithm | 9 |
| 3.14. Maximum Density Subgraph | 9 |
| 3.15. Maximum-Weight Closure | 9 |
| 3.16. Maximum Weighted Independent Set in a Bipartite Graph | 9 |
| 3.17. Synchronizing word problem | 10 |
| 3.18. Max flow with lower bounds on edges | 10 |
| 3.19. Tutte matrix for general matching | 10 |
| 3.20. Heavy Light Decomposition | 10 |
| 3.21. Centroid Decomposition | 10 |
| 3.22. Least Common Ancestor | 10 |
| 3.23. Counting Spanning Trees | 11 |
| 3.24. Erdős-Gallai Theorem | 11 |
| 3.25. Tree Isomorphism | 11 |
| 4. Strings | 11 |
| 4.1. Knuth-Morris-Pratt | 11 |
| 4.2. Suffix Array | 11 |
| 4.3. Longest Common Prefix | 11 |
| 4.4. Aho-Corasick Trie | 11 |
| 4.5. Palindromic Tree | 12 |
| 4.6. Z Algorithm | 12 |
| 4.7. Booth's Minimum String Rotation | 12 |
| 4.8. Hashing | 12 |
| 5. Dynamic Programming | 13 |
| 5.1. Longest Common Subsequence | 13 |
| 5.2. Longest Increasing Subsequence | 13 |
| 5.3. Traveling Salesman | 13 |
| 6. Number Theory | 13 |

| | |
|---|----|
| 6.1. Eratosthenes Prime Sieve | 13 |
| 6.2. Divisor Sieve | 13 |
| 6.3. Number/Sum of Divisors | 13 |
| 6.4. Möbius Sieve | 13 |
| 6.5. Möbius Inversion | 13 |
| 6.6. GCD Subset Counting | 13 |
| 6.7. Euler Totient | 13 |
| 6.8. Euler Phi Sieve | 13 |
| 6.9. Extended Euclidean | 13 |
| 6.10. Modular Inverse | 13 |
| 6.11. Modulo Solver | 13 |
| 6.12. Linear Diophantine | 13 |
| 6.13. Chinese Remainder Theorem | 13 |
| 7. Algebra | 14 |
| 7.1. Fast Fourier Transform | 14 |
| 7.2. FFT Polynomial Multiplication | 14 |
| 7.3. Polynomial Long Division | 14 |
| 7.4. Matrix Multiplication | 14 |
| 7.5. Matrix Power | 14 |
| 7.6. Fibonacci Matrix | 14 |
| 7.7. Gauss-Jordan/Matrix Determinant | 14 |
| 8. Combinatorics | 14 |
| 8.1. Lucas Theorem | 14 |
| 8.2. Granville's Theorem | 14 |
| 8.3. Derangements | 15 |
| 8.4. Factoradics | 15 |
| 8.5. k th Permutation | 15 |
| 8.6. Catalan Numbers | 15 |
| 8.7. Stirling Numbers | 15 |
| 8.8. Partition Function | 15 |
| 9. Geometry | 15 |
| 9.1. Dots and Cross Products | 15 |
| 9.2. Angles and Rotations | 15 |
| 9.3. Spherical Coordinates | 15 |
| 9.4. Point Projection | 15 |
| 9.5. Great Circle Distance | 15 |
| 9.6. Point/Line/Plane Distances | 16 |
| 9.7. Intersections | 16 |
| 9.8. Polygon Areas | 16 |
| 9.9. Polygon Centroid | 16 |
| 9.10. Convex Hull | 16 |
| 9.11. Point in Polygon | 16 |
| 9.12. Cut Polygon by a Line | 17 |
| 9.13. Triangle Centers | 17 |
| 9.14. Convex Polygon Intersection | 17 |
| 9.15. Pick's Theorem for Lattice Points | 17 |
| 9.16. Minimum Enclosing Circle | 17 |
| 9.17. Shamos Algorithm | 17 |
| 9.18. k D Tree | 17 |
| 9.19. Line Sweep (Closest Pair) | 17 |
| 9.20. Line upper/lower envelope | 18 |
| 9.21. Formulas | 18 |
| 10. Other Algorithms | 18 |
| 10.1. Coordinate Compression | 18 |
| 10.2. 2SAT | 18 |
| 10.3. Nth Permutation | 18 |

| | |
|--|----|
| 10.4. Floyd's Cycle-Finding | 18 |
| 10.5. Simulated Annealing | 18 |
| 10.6. Hexagonal Grid Algorithms | 18 |
| 11. Useful Information (CLEAN THIS UP!!) | 19 |
| 12. Misc | 19 |
| 12.1. Debugging Tips | 19 |
| 12.2. Solution Ideas | 19 |
| 13. Formulas | 20 |
| 13.1. Physics | 20 |
| 13.2. Markov Chains | 20 |
| 13.3. Burnside's Lemma | 20 |
| 13.4. Bézout's identity | 20 |
| 13.5. Misc | 20 |

1. CODE TEMPLATES

```
#include <bits/stdc++.h> -----//001
typedef long long ll; -----//002
typedef unsigned long long ull; -----//003
typedef std::pair<int, int> ii; -----//004
typedef std::pair<int, ii> iii; -----//005
typedef std::vector<int> vi; -----//006
typedef std::vector<vi> vvi; -----//007
typedef std::vector<ii> vii; -----//008
typedef std::vector<iii> viii; -----//009
const int INF = ~(1<<31); -----//00a
const ll LINF = (1LL << 60); -----//00b
const int MAXN = 1e5+1; -----//00c
const double EPS = 1e-9; -----//00d
const double pi = acos(-1); -----//00e
```

2. DATA STRUCTURES

```
2.1. Union Find.
struct union_find { -----//2f3
- vi p; union_find(int n) : p(n, -1) { } -----//2f4
- int find(int x) { return p[x] < 0 ? x : p[x] = find(p[x]); }
- bool unite(int x, int y) { -----//2f6
--- int xp = find(x), yp = find(y); -----//2f7
--- if (xp == yp) return false; -----//2f8
--- if (p[xp] > p[yp]) std::swap(xp,yp); -----//2f9
--- p[xp] += p[yp], p[yp] = xp; -----//2fa
--- return true; -----//2fb
- } -----//2fc
- int size(int x) { return -p[find(x)]; } -----//2fd
}; -----//2fe

2.2. Segment Tree.
2.2.1. Recursive, Point-update Segment Tree.
struct segtree { -----//1df
- int i, j, val; -----//1e0
- segtree *l, *r; -----//1e1
- segtree(vi &ar, int _i, int _j) : i(_i), j(_j) { -----//1e2
--- if (i == j) { -----//1e3
---- val = ar[i]; -----//1e4
```

```

- l = r = NULL; -----//1e5 struct segtree { -----//1af build(ar, 1, 0, n-1); -----//130
- } else { -----//1e6 - int i, j, val, temp_val = 0; -----//1b0 - } -----//131
- int k = (i+j) >> 1; -----//1e7 - segtree *l, *r; -----//1b1 - void build(vi &ar, int p, int i, int j) { -----//132
- l = new segtree(ar, i, k); -----//1e8 - segtree(vi &ar, int _i, int _j) : i(_i), j(_j) { -----//1b2 - deltas[p] = 0; -----//133
- r = new segtree(ar, k+1, j); -----//1e9 - if (i == j) { -----//1b3 - if (i == j) -----//134
- val = l->val + r->val; -----//1ea - val = ar[i]; -----//1b4 - vals[p] = ar[i]; -----//135
- } -----//1eb - l = r = NULL; -----//1b5 - else { -----//136
- } -----//1ec - } else { -----//1b6 - int k = (i + j) / 2; -----//137
- void update(int _i, int _val) { -----//1ed - int k = (i + j) >> 1; -----//1b7 - build(ar, p<<1, i, k); -----//138
- if (_i <= i and j <= _i) { -----//1ee - l = new segtree(ar, i, k); -----//1b8 - build(ar, p<<1|1, k+1, j); -----//139
- val += _val; -----//1ef - r = new segtree(ar, k+1, j); -----//1b9 - pull(p); -----//13a
- } else if (_i < i or j < _i) { -----//1f0 - val = l->val + r->val; -----//1ba - } -----//13b
- // do nothing -----//1f1 - } -----//1bb - } -----//13c
- } else { -----//1f2 - } -----//1bc - void pull(int p) { -----//13d
- l->update(_i, _val); -----//1f3 - void visit() { -----//1bd - vals[p] = vals[p<<1] + vals[p<<1|1]; -----//13e
- r->update(_i, _val); -----//1fa - if (temp_val) { -----//1be - } -----//13f
- val = l->val + r->val; -----//1fb - val += (j-i+1) * temp_val; -----//1bf - void push(int p, int i, int j) { -----//140
- } -----//1fc - if (!l) { -----//1c0 - if (deltas[p]) { -----//141
- } -----//1fd - l->temp_val += temp_val; -----//1c1 - vals[p] += (j - i + 1) * deltas[p]; -----//142
- int query(int _i, int _j) { -----//1fe - r->temp_val += temp_val; -----//1c2 - if (i != j) { -----//143
- if (_i <= i and j <= _j) { -----//1ff - } -----//1c3 - deltas[p<<1] += deltas[p]; -----//144
- return val; -----//1fa - temp_val = 0; -----//1c4 - deltas[p<<1|1] += deltas[p]; -----//145
- } else if (_j < i or j < _i) { -----//1fb - } -----//1c5 - } -----//146
- return 0; -----//1fc - } -----//1c6 - deltas[p] = 0; -----//147
- } else { -----//1fd - void increase(int _i, int _j, int _inc) { -----//1c7 - } -----//148
- return l->query(_i, _j) + r->query(_i, _j); -----//1fe - visit(); -----//1c8 - } -----//149
- } -----//1ff - if (_i <= i && j <= _j) { -----//1c9 - void update(int _i, int _j, int v, -----//14a
- } -----//200 - temp_val += _inc; -----//1ca - int p, int i, int j) { -----//14b
- }; -----//201 - visit(); -----//1cb - push(p, i, j); -----//14c
- -----// do nothing -----//1cc - if (_i <= i && j <= _j) { -----//14d
- } else { -----//1cd - deltas[p] += v; -----//14e
- l->increase(_i, _j, _inc); -----//1ce - push(p, i, j); -----//14f
- r->increase(_i, _j, _inc); -----//1cf - } else if (_j < i || j < _i) { -----//150
- val = l->val + r->val; -----//1d0 - // do nothing -----//151
- } -----//1d1 - } else { -----//152
- } -----//1d2 - int k = (i + j) / 2; -----//153
- } -----//1d3 - update(_i, _j, v, p<<1, i, k); -----//154
- int query(int _i, int _j) { -----//1d4 - update(_i, _j, v, p<<1|1, k+1, j); -----//155
- visit(); -----//1d5 - pull(p); -----//156
- if (_i <= i and j <= _j) { -----//1d6 - } -----//157
- return val; -----//1d7 - } -----//158
- } else if (_j < i || j < _i) { -----//1d8 - int query(int _i, int _j, -----//159
- return 0; -----//1d9 - int p, int i, int j) { -----//15a
- } else { -----//1da - push(p, i, j); -----//15b
- return l->query(_i, _j) + r->query(_i, _j); -----//1db - if (_i <= i and j <= _j) { -----//15c
- } -----//1dc - return vals[p]; -----//15d
- } -----//1dd - } else if (_j < i || j < _i) { -----//15e
- } -----//1de - return 0; -----//15f
- }; -----//1df - } else { -----//160
- -----//202 - int k = (i + j) / 2; -----//161
- -----//203 - return query(_i, _j, p<<1, i, k) + -----//162
- -----//204 - query(_i, _j, p<<1|1, k+1, j); -----//163
- -----//205 - } -----//164
- -----//206 - } -----//165
- -----//207 - }; -----//166

```

2.2.2. Iterative, Point-update Segment Tree.

```

struct segtree { -----//167
- int n; -----//168
- int *vals; -----//169
- segtree(vi &ar, int n) { -----//16a
- this->n = n; -----//16b
- vals = new int[2*n]; -----//16c
- for (int i = 0; i < n; ++i) -----//16d
- vals[i+n] = ar[i]; -----//16e
- for (int i = n-1; i > 0; --i) -----//16f
- vals[i] = vals[i<<1] + vals[i<<1|1]; -----//170
- } -----//171
- void update(int i, int v) { -----//172
- for (vals[i += n] += v; i > 1; i >>= 1) -----//173
- vals[i>>1] = vals[i] + vals[i^1]; -----//174
- } -----//175
- int query(int l, int r) { -----//176
- int res = 0; -----//177
- for (l += n, r += n+1; l < r; l >>= 1, r >>= 1) { -----//178
- if (l&1) res += vals[l++]; -----//179
- if (r&1) res += vals[--r]; -----//17a
- } -----//17b
- return res; -----//17c
- } -----//17d
}; -----//17e

```

2.2.4. Array-based, Range-update Segment Tree.

```

struct segtree { -----//12a
- int n, *vals, *deltas; -----//12b
- segtree(vi &ar) { -----//12c
- n = ar.size(); -----//12d
- vals = new int[4*n]; -----//12e
- deltas = new int[4*n]; -----//12f
}

```

2.2.3. Pointer-based, Range-update Segment Tree.

2.2.5. Persistent Segment Tree (Point-update).

```
struct node { int l, r, lid, rid, val; }; //17f
struct segtree { //180
- node *nodes; //181
- int n, node_cnt = 0; //182
- segtree(int n, int capacity) { //183
-- this->n = n; //184
-- nodes = new node[capacity]; //185
- } //186
- int build (vi &ar, int l, int r) { //187
-- if (l > r) return -1; //188
-- int id = node_cnt++; //189
-- nodes[id].l = l; //18a
-- nodes[id].r = r; //18b
-- if (l == r) { //18c
-- nodes[id].lid = -1; //18d
-- nodes[id].rid = -1; //18e
-- nodes[id].val = ar[l]; //18f
- } else { //190
-- int m = (l + r) / 2; //191
-- nodes[id].lid = build(ar, l, m); //192
-- nodes[id].rid = build(ar, m+1, r); //193
-- nodes[id].val = nodes[nodes[id].lid].val + //194
-- nodes[nodes[id].rid].val; //195
- } //196
- return id; //197
- } //198
- int update(int id, int idx, int delta) { //199
-- if (id == -1) //19a
-- return -1; //19b
-- if (idx < nodes[id].l or nodes[id].r < idx) //19c
-- return id; //19d
-- int nid = node_cnt++; //19e
-- nodes[nid].l = nodes[id].l; //19f
-- nodes[nid].r = nodes[id].r; //1a0
-- nodes[nid].lid = update(nodes[id].lid, idx, delta); //1a1
-- nodes[nid].rid = update(nodes[id].rid, idx, delta); //1a2
-- nodes[nid].val = nodes[id].val + delta; //1a3
-- return nid; //1a4
- } //1a5
- int query(int id, int l, int r) { //1a6
-- if (r < nodes[id].l or nodes[id].r < l) //1a7
-- return 0; //1a8
-- if (l <= nodes[id].l and nodes[id].r <= r) //1a9
-- return nodes[id].val; //1aa
-- return query(nodes[id].lid, l, r) + //1ab
-- query(nodes[id].rid, l, r); //1ac
- } //1ad
}; //1ae

//17f ar[i] = new int[m]; //113
//180 for (int j = 0; j < m; ++j) //114
//181 ar[i][j] = 0; //115
//182 } //116
//183 } //117
//184 void update(int x, int y, int v) { //118
//185 ar[x + n][y + m] = v; //119
//186 for (int i = x + n; i > 0; i >= 1) { //11a
//187 for (int j = y + m; j > 0; j >= 1) { //11b
//188 ar[i>>1][j] = min(ar[i][j], ar[i^1][j]); //11c
//189 ar[i][j>>1] = min(ar[i][j], ar[i][j^1]); //11d
//18a }}} // just call update one by one to build //11e
//18b int query(int x1, int x2, int y1, int y2) { //11f
//18c int s = INF; //120
//18d if(~x2) for(int a=x1+n, b=x2+n+1; a<b; a>=1, b>=1) { //121
//18e if (a & 1) s = min(s, query(a++, -1, y1, y2)); //122
//18f if (b & 1) s = min(s, query(--b, -1, y1, y2)); //123
//190 } else for (int a=y1+m, b=y2+m+1; a<b; a>=1, b>=1) { //124
//191 if (a & 1) s = min(s, ar[x1][a++]); //125
//192 if (b & 1) s = min(s, ar[x1][--b]); //126
//193 } return s; //127
//194 } //128
//195 }; //129

2.3. Fenwick Tree.
2.3.1. Fenwick Tree w/ Point Queries.
struct fenwick { //0c7
- vi ar; //0c8
- fenwick(vi &ar) : ar(_ar.size(), 0) { //0c9
-- for (int i = 0; i < ar.size(); ++i) { //0ca
-- ar[i] += _ar[i]; //0cb
-- int j = i | (i+1); //0cc
-- if (j < ar.size()) //0cd
-- ar[j] += ar[i]; //0ce
-- } //0cf
- } //0d0
- int sum(int i) { //0d1
-- int res = 0; //0d2
-- for (; i >= 0; i = (i & (i+1)) - 1) //0d3
-- res += ar[i]; //0d4
-- return res; //0d5
- } //0d6
- int sum(int i, int j) { return sum(j) - sum(i-1); } //0d7
- void add(int i, int val) { //0d8
-- for (; i < ar.size(); i |= i+1) //0d9
-- ar[i] += val; //0da
- } //0db
- int get(int i) { //0dc
-- int res = ar[i]; //0dd
-- if (i) { //0de
-- int lca = (i & (i+1)) - 1; //0df
-- for (--i; i != lca; i = (i&(i+1))-1) //0e0
-- res -= ar[i]; //0e1
-- } //0e2
-- return res; //0e3
- } //0e4

- void set(int i, int val) { add(i, -get(i) + val); } //0e5
- // range update, point query // //0e6
- void add(int i, int j, int val) { //0e7
-- add(i, val); //0e8
-- add(j+1, -val); //0e9
- } //0ea
- int get1(int i) { return sum(i); } //0eb
- ////////////////////////////////////////////////// //0ec
}; //0ed

2.3.2. Fenwick Tree w/ Max Queries.
struct fenwick { //0ee
- vi ar; //0ef
- fenwick(vi &ar) : ar(_ar.size(), 0) { //0f0
-- for (int i = 0; i < ar.size(); ++i) { //0f1
-- ar[i] = std::max(ar[i], _ar[i]); //0f2
-- int j = i | (i+1); //0f3
-- if (j < ar.size()) //0f4
-- ar[j] = std::max(ar[j], ar[i]); //0f5
-- } //0f6
- } //0f7
- void set(int i, int v) { //0f8
-- for (; i < ar.size(); i |= i+1) //0f9
-- ar[i] = std::max(ar[i], v); //0fa
- } //0fb
- // max[0..i] //0fc
- int max(int i) { //0fd
-- int res = -INF; //0fe
-- for (; i >= 0; i = (i & (i+1)) - 1) //0ff
-- res = std::max(res, ar[i]); //100
-- return res; //101
- } //102
}; //103

2.4. Treap.
2.4.1. Explicit Treap.
2.4.2. Implicit Treap.
struct cartree { //28b
- typedef struct _Node { //28c
-- int node_val, subtree_val, delta, prio, size; //28d
-- _Node *l, *r; //28e
-- _Node(int val) : node_val(val), subtree_val(val), //28f
-- delta(0), prio((rand()<16)^rand()), size(1), //290
-- l(NULL), r(NULL) {} //291
-- ~_Node() { delete l; delete r; } //292
- } *Node; //293
- int get_subtree_val(Node v) { //294
-- return v ? v->subtree_val : 0; } //295
- int get_size(Node v) { return v ? v->size : 0; } //296
- void apply_delta(Node v, int delta) { //297
-- if (!v) return; //298
-- v->delta += delta; //299
-- v->node_val += delta; //29a
-- v->subtree_val += delta * get_size(v); //29b
- } //29c
- void push_delta(Node v) { //29d
```

```
-- if (!v) return; -----//29e -- insert(new _Node(val), key); -----//2d6 - } // push down lazy flags to children (editable) -----//24e
-- apply_delta(v->l, v->delta); -----//29f - } -----//2d7 - void push(node *p) { -----//24f
-- apply_delta(v->r, v->delta); -----//2a0 - void erase(int key) { -----//2d8 - if (p != null && p->reverse) { -----//250
-- v->delta = 0; -----//2a1 - Node l, m, r; -----//2d9 - swap(p->left, p->right); -----//251
- } -----//2a2 - split(root, key + 1, m, r); -----//2da - p->left->reverse ^= 1; -----//252
- void update(Node v) { -----//2a3 - split(m, key, l, m); -----//2db - p->right->reverse ^= 1; -----//253
-- if (!v) return; -----//2a4 - delete m; -----//2dc - p->reverse ^= 1; -----//254
-- v->subtree_val = get_subtree_val(v->l) + v->node_val -----//2a5 - root = merge(l, r); -----//2dd - } } // assign son to be the new child of p -----//255
-- + get_subtree_val(v->r); -----//2a6 - } -----//2de - void link(node *p, node *son, int d) { -----//256
-- v->size = get_size(v->l) + 1 + get_size(v->r); -----//2a7 - int query(int a, int b) { -----//2df - p->get(d) = son; -----//257
- } -----//2a8 - Node l1, r1; -----//2e0 - son->parent = p; } -----//258
- Node merge(Node l, Node r) { -----//2a9 - split(root, b+1, l1, r1); -----//2e1 - int dir(node *p, node *son) { -----//259
-- push_delta(l); push_delta(r); -----//2aa - Node l2, r2; -----//2e2 - return p->left == son ? 0 : 1; } -----//25a
-- if (!l || !r) return l ? l : r; -----//2ab - split(l1, a, l2, r2); -----//2e3 - void rotate(node *x, int d) { -----//25b
-- if (l->size <= r->size) { -----//2ac - int res = get_subtree_val(r2); -----//2e4 - node *y = x->get(d), *z = x->parent; -----//25c
-- l->r = merge(l->r, r); -----//2ad - l1 = merge(l2, r2); -----//2e5 - link(x, y->get(d ^ 1), d); -----//25d
-- update(l); -----//2ae - root = merge(l1, r1); -----//2e6 - link(y, x, d ^ 1); -----//25e
-- return l; -----//2af - return res; -----//2e7 - link(z, y, dir(z, x)); -----//25f
-- } else { -----//2b0 - } -----//2e8 - pull(x); pull(y); } -----//260
-- r->l = merge(l, r->l); -----//2b1 - void update(int a, int b, int delta) { -----//2e9 - node* splay(node *p) { // splay node p to root -----//261
-- update(r); -----//2b2 - Node l1, r1; -----//2ea - while (p->parent != null) { -----//262
-- return r; -----//2b3 - split(root, b+1, l1, r1); -----//2eb - node *m = p->parent, *g = m->parent; -----//263
- } -----//2b4 - Node l2, r2; -----//2ec - push(g); push(m); push(p); -----//264
- } -----//2b5 - split(l1, a, l2, r2); -----//2ed - int dm = dir(m, p), dg = dir(g, m); -----//265
- void split(Node v, int key, Node &l, Node &r) { -----//2b6 - apply_delta(r2, delta); -----//2ee - if (g == null) rotate(m, dm); -----//266
-- push_delta(v); -----//2b7 - l1 = merge(l2, r2); -----//2ef - else if (dm == dg) rotate(g, dg), rotate(m, dm); -----//267
-- l = r = NULL; -----//2b8 - root = merge(l1, r1); -----//2f0 - else rotate(m, dm), rotate(g, dg); -----//268
-- if (!v) return; -----//2b9 - } -----//2f1 - } return root = p; } -----//269
-- if (key <= get_size(v->l)) { -----//2ba - int size() { return get_size(root); } }; -----//2f2
-- split(v->l, key, l, v->l); -----//2bb
-- r = v; -----//2bc
-- } else { -----//2bd
-- split(v->r, key - get_size(v->l) - 1, v->r, r); -----//2be
-- l = v; -----//2bf
- } -----//2c0
-- update(v); -----//2c1
- } -----//2c2
- Node root; -----//2c3
public: -----//2c4
- cartree() : root(NULL) {} -----//2c5
- ~cartree() { delete root; } -----//2c6
- int get(Node v, int key) { -----//2c7
-- push_delta(v); -----//2c8
-- if (key < get_size(v->l)) -----//2c9
-- return get(v->l, key); -----//2ca
-- else if (key > get_size(v->l)) -----//2cb
-- return get(v->r, key - get_size(v->l) - 1); -----//2cc
-- return v->node_val; -----//2cd
- } -----//2ce
- int get(int key) { return get(root, key); } -----//2cf
- void insert(Node item, int key) { -----//2d0
-- Node l, r; -----//2d1
-- split(root, key, l, r); -----//2d2
-- root = merge(merge(l, item), r); -----//2d3
- } -----//2d4
- void insert(int key, int val) { -----//2d5
-- insert(new _Node(val), key); -----//2d6
-- void erase(int key) { -----//2d8
-- Node l, m, r; -----//2d9
-- split(root, key + 1, m, r); -----//2da
-- split(m, key, l, m); -----//2db
-- delete m; -----//2dc
-- root = merge(l, r); -----//2dd
- } -----//2de
-- int query(int a, int b) { -----//2df
-- Node l1, r1; -----//2e0
-- split(root, b+1, l1, r1); -----//2e1
-- Node l2, r2; -----//2e2
-- split(l1, a, l2, r2); -----//2e3
-- int res = get_subtree_val(r2); -----//2e4
-- l1 = merge(l2, r2); -----//2e5
-- root = merge(l1, r1); -----//2e6
-- return res; -----//2e7
- } -----//2e8
-- void update(int a, int b, int delta) { -----//2e9
-- Node l1, r1; -----//2ea
-- split(root, b+1, l1, r1); -----//2eb
-- Node l2, r2; -----//2ec
-- split(l1, a, l2, r2); -----//2ed
-- apply_delta(r2, delta); -----//2ee
-- l1 = merge(l2, r2); -----//2ef
-- root = merge(l1, r1); -----//2f0
- } -----//2f1
-- int size() { return get_size(root); } }; -----//2f2

2.4.3. Persistent Treap.

2.5. Splay Tree.
struct node *null; -----//236
struct node { -----//237
- node *left, *right, *parent; -----//238
- bool reverse; int size, value; -----//239
- node*& get(int d) {return d == 0 ? left : right;} -----//23a
- node(int v=0): reverse(0), size(0), value(v) { -----//23b
- left = right = parent = null ? null : this; -----//23c
- }; -----//23d
} -----//23e
struct SplayTree { -----//23e
- node *root; -----//23f
- SplayTree(int arr[] = NULL, int n = 0) { -----//240
-- if (!null) null = new node(); -----//241
-- root = build(arr, n); -----//242
- } // build a splay tree based on array values -----//243
- node* build(int arr[], int n) { -----//244
-- if (n == 0) return null; -----//245
-- int mid = n >> 1; -----//246
-- node *p = new node(arr ? arr[mid] : 0); -----//247
-- link(p, build(arr, mid), 0); -----//248
-- link(p, build(arr? arr+mid+1 : NULL, n-mid-1), 1); -----//249
-- pull(p); return p; -----//24a
- } // pull information from children (editable) -----//24b
- void pull(node *p) { -----//24c
-- p->size = p->left->size + p->right->size + 1; -----//24d
-- }
-- } // push down lazy flags to children (editable) -----//24e
-- void push(node *p) { -----//24f
-- if (p != null && p->reverse) { -----//250
-- swap(p->left, p->right); -----//251
-- p->left->reverse ^= 1; -----//252
-- p->right->reverse ^= 1; -----//253
-- p->reverse ^= 1; -----//254
-- } } // assign son to be the new child of p -----//255
-- void link(node *p, node *son, int d) { -----//256
-- p->get(d) = son; -----//257
-- son->parent = p; } -----//258
-- int dir(node *p, node *son) { -----//259
-- return p->left == son ? 0 : 1; } -----//25a
-- void rotate(node *x, int d) { -----//25b
-- node *y = x->get(d), *z = x->parent; -----//25c
-- link(x, y->get(d ^ 1), d); -----//25d
-- link(y, x, d ^ 1); -----//25e
-- link(z, y, dir(z, x)); -----//25f
-- pull(x); pull(y); } -----//260
-- node* splay(node *p) { // splay node p to root -----//261
-- while (p->parent != null) { -----//262
-- node *m = p->parent, *g = m->parent; -----//263
-- push(g); push(m); push(p); -----//264
-- int dm = dir(m, p), dg = dir(g, m); -----//265
-- if (g == null) rotate(m, dm); -----//266
-- else if (dm == dg) rotate(g, dg), rotate(m, dm); -----//267
-- else rotate(m, dm), rotate(g, dg); -----//268
-- } return root = p; } -----//269
-- node* get(int k) { // get the node at index k -----//26a
-- node *p = root; -----//26b
-- while (push(p), p->left->size != k) { -----//26c
-- if (k < p->left->size) p = p->left; -----//26d
-- else k -= p->left->size + 1, p = p->right; -----//26e
-- } -----//26f
-- return p == null ? null : splay(p); -----//270
- } // keep the first k nodes, the rest in r -----//271
-- void split(node *r, int k) { -----//272
-- if (k == 0) {r = root; root = null; return;} -----//273
-- r = get(k - 1)->right; -----//274
-- root->right = r->parent = null; -----//275
-- pull(root); } -----//276
-- void merge(node *r) { //merge current tree with r -----//277
-- if (root == null) {root = r; return;} -----//278
-- link(get(root->size - 1), r, 1); -----//279
-- pull(root); } -----//27a
-- void assign(int k, int val) { // assign arr[k]= val -----//27b
-- get(k)->value = val; pull(root); } -----//27c
-- void reverse(int L, int R) { // reverse arr[L...R] -----//27d
-- node *m, *r; split(r, R + 1); split(m, L); -----//27e
-- m->reverse ^= 1; push(m); merge(m); merge(r); -----//27f
- } // insert a new node before the node at index k -----//280
-- node* insert(int k, int v) { -----//281
-- node *r; split(r, k); -----//282
-- node *p = new node(v); p->size = 1; -----//283
-- link(root, p, 1); merge(r); -----//284
-- return p; } -----//285
```

```
- void erase(int k) { // erase node at index k -----//286
-- node *r, *m; -----//287
-- split(r, k + 1); split(m, k); -----//288
-- merge(r); delete m;} -----//289
}; -----//28a
```

2.6. Ordered Statistics Tree.

```
#include <ext/pb_ds/assoc_container.hpp> -----//104
#include <ext/pb_ds/tree_policy.hpp> -----//105
using namespace __gnu_pbds; -----//106
template <typename T> -----//107
using indexed_set = std::tree<T, null_type, less<T>, -----//108
splay_tree_tag, tree_order_statistics_node_update>; -----//109
// indexed_set<int> t; t.insert(...); -----//10a
// t.find_by_order(index); // 0-based -----//10b
// t.order_of_key(key); -----//10c
```

2.7. Sparse Table.

2.7.1. 1D Sparse Table.

```
int lg[MAXN+1], spt[20][MAXN]; -----//202
void build(vi &arr, int n) { -----//203
- for (int i = 2; i <= n; ++i) lg[i] = lg[i>>1] + 1; -----//204
- for (int i = 0; i < n; ++i) spt[0][i] = arr[i]; -----//205
- for (int j = 0; (2 << j) <= n; ++j) -----//206
-- for (int i = 0; i + (2 << j) <= n; ++i) -----//207
--- spt[j+1][i] = std::min(spt[j][i], spt[j][i+(1<<j)]);//208
} -----//209
int query(int a, int b) { -----//20a
- int k = lg[b-a+1], ab = b - (1<<k) + 1; -----//20b
- return std::min(spt[k][a], spt[k][ab]); -----//20c
} -----//20d
```

2.7.2. 2D Sparse Table.

```
const int N = 100, LGN = 20; -----//20e
int lg[N], A[N][N], st[LGN][LGN][N][N]; -----//20f
void build(int n, int m) { -----//210
- for(int k=2; k<=std::max(n,m); ++k) lg[k] = lg[k>>1]+1; //211
- for(int i = 0; i < n; ++i) -----//212
-- for(int j = 0; j < m; ++j) -----//213
--- st[0][0][i][j] = A[i][j]; -----//214
- for(int bj = 0; (2 << bj) <= m; ++bj) -----//215
-- for(int j = 0; j + (2 << bj) <= m; ++j) -----//216
--- for(int i = 0; i < n; ++i) -----//217
---- st[0][bj+1][i][j] = -----//218
----- std::max(st[0][bj][i][j], -----//219
----- st[0][bj][i][j + (1 << bj)]); -----//21a
- for(int bi = 0; (2 << bi) <= n; ++bi) -----//21b
-- for(int i = 0; i + (2 << bi) <= n; ++i) -----//21c
--- for(int j = 0; j < m; ++j) -----//21d
---- st[bi+1][0][i][j] = -----//21e
----- std::max(st[bi][0][i][j], -----//21f
----- st[bi][0][i + (1 << bi)][j]); -----//220
- for(int bi = 0; (2 << bi) <= n; ++bi) -----//221
-- for(int i = 0; i + (2 << bi) <= n; ++i) -----//222
--- for(int bj = 0; (2 << bj) <= m; ++bj) -----//223
---- for(int j = 0; j + (2 << bj) <= m; ++j) { -----//224
----- int ik = i + (1 << bi); -----//225
```

```
int jk = j + (1 << bj); -----//226
st[bi+1][bj+1][i][j] = -----//227
std::max(std::max(st[bi][bj][i][j], -----//228
st[bi][bj][ik][j]), -----//229
std::max(st[bi][bj][i][jk], -----//22a
st[bi][bj][ik][jk])); -----//22b
} -----//22c
} -----//22d
int query(int x1, int x2, int y1, int y2) { -----//22e
- int kx = lg[x2 - x1 + 1], ky = lg[y2 - y1 + 1]; -----//22f
- int x12 = x2 - (1<<kx) + 1, y12 = y2 - (1<<ky) + 1; -----//230
- return std::max(std::max(st[kx][ky][x1][y1], -----//231
st[kx][ky][x1][y12]), -----//232
std::max(st[kx][ky][x12][y1], -----//233
st[kx][ky][x12][y12])); -----//234
} -----//235
```

3. GRAPHS

Using adjacency list:

```
struct graph { -----//542
- int n, *dist; -----//543
- vii *adj; -----//544
- graph(int n) { -----//545
-- this->n = n; -----//546
-- adj = new vii[n]; -----//547
-- dist = new int[n]; -----//548
- } -----//549
- void add_edge(int u, int v, int w) { -----//54a
-- adj[u].push_back({v, w}); -----//54b
-- // adj[v].push_back({u, w}); -----//54c
- } -----//54d
}; -----//54e
```

Using adjacency matrix:

```
struct graph { -----//54f
- int n, **mat; -----//550
- graph(int n) { -----//551
-- this->n = n; -----//552
-- mat = new int*[n]; -----//553
-- for (int i = 0; i < n; ++i) { -----//554
-- mat[i] = new int[n]; -----//555
-- for (int j = 0; j < n; ++j) -----//556
-- mat[i][j] = INF; -----//557
-- mat[i][i] = 0; -----//558
- } -----//559
- } -----//55a
- void add_edge(int u, int v, int w) { -----//55b
-- mat[u][v] = std::min(mat[u][v], w); -----//55c
-- // mat[v][u] = std::min(mat[v][u], w); -----//55d
- } -----//55e
}; -----//55f
```

Using edge list:

```
struct graph { -----//560
- int n; -----//561
- std::vector<iii> edges; -----//562
- graph(int n) : n(n) {} -----//563
```

```
- void add_edge(int u, int v, int w) { -----//564
-- edges.push_back({w, {u, v}}); -----//565
- } -----//566
}; -----//567
```

3.1. Single-Source Shortest Paths.

3.1.1. Dijkstra.

```
#include "graph_template_adjlist.cpp" -----//7a1
// insert inside graph; needs n, dist[], and adj[] -----//7a2
void dijkstra(int s) { -----//7a3
- for (int u = 0; u < n; ++u) -----//7a4
-- dist[u] = INF; -----//7a5
- dist[s] = 0; -----//7a6
- std::priority_queue<ii, vii, std::greater<ii> > pq; -----//7a7
- pq.push({0, s}); -----//7a8
- while (!pq.empty()) { -----//7a9
-- int u = pq.top().second; -----//7aa
-- int d = pq.top().first; -----//7ab
-- pq.pop(); -----//7ac
-- if (dist[u] < d) -----//7ad
-- continue; -----//7ae
-- dist[u] = d; -----//7af
-- for (auto &e : adj[u]) { -----//7b0
-- int v = e.first; -----//7b1
-- int w = e.second; -----//7b2
-- if (dist[v] > dist[u] + w) { -----//7b3
-- dist[v] = dist[u] + w; -----//7b4
-- pq.push({dist[v], v}); -----//7b5
-- } -----//7b6
-- } -----//7b7
- } -----//7b8
} -----//7b9
```

3.1.2. Bellman-Ford.

```
#include "graph_template_adjlist.cpp" -----//78d
// insert inside graph; needs n, dist[], and adj[] -----//78e
void bellman_ford(int s) { -----//78f
- for (int u = 0; u < n; ++u) -----//790
-- dist[u] = INF; -----//791
- dist[s] = 0; -----//792
- for (int i = 0; i < n-1; ++i) -----//793
-- for (int u = 0; u < n; ++u) -----//794
--- for (auto &e : adj[u]) -----//795
---- if (dist[u] + e.second < dist[e.first]) -----//796
----- dist[e.first] = dist[u] + e.second; -----//797
} -----//798
// you can call this after running bellman_ford() -----//799
bool has_neg_cycle() { -----//79a
- for (int u = 0; u < n; ++u) -----//79b
-- for (auto &e : adj[u]) -----//79c
--- if (dist[e.first] > dist[u] + e.second) -----//79d
-- return true; -----//79e
- return false; -----//79f
} -----//7a0
```


3.1.3. SPFA.

```
struct edge { -----//7c3
- int v; long long cost; -----//7c4
- edge(int v, long long cost): v(v), cost(cost) {} -----//7c5
}; -----//7c6
long long dist[N]; int vis[N]; bool inq[N]; -----//7c7
void spfa(vector<edge*> adj[], int n, int s) { -----//7c8
- fill(dist, dist + n, LLONG_MAX); -----//7c9
- fill(vis, vis + n, 0); -----//7ca
- fill(inq, inq + n, false); -----//7cb
- queue<int> q; q.push(s); -----//7cc
- for (dist[s] = 0; !q.empty(); q.pop()) { -----//7cd
-- int u = q.front(); inq[u] = false; -----//7ce
-- if (++vis[u] >= n) dist[u] = LLONG_MIN; -----//7cf
-- for (int i = 0; i < adj[u].size(); ++i) { -----//7d0
--- edge& e = *adj[u][i]; -----//7d1
--- // uncomment below for min cost max flow -----//7d2
--- // if (e.cap <= e.flow) continue; -----//7d3
--- int v = e.v; -----//7d4
--- long long w = vis[u] >= n ? 0LL : e.cost; -----//7d5
--- if (dist[u] + w < dist[v]) { -----//7d6
---- dist[v] = dist[u] + w; -----//7d7
---- if (!inq[v]) { -----//7d8
----- inq[v] = true; -----//7d9
----- q.push(v); -----//7da
----- }}}} -----//7db
```

3.2. All-Pairs Shortest Paths.

3.2.1. Floyd-Washall.

```
#include "graph_template_adjmat.cpp" -----//7ba
// insert inside graph; needs n and mat[][] -----//7bb
void floyd_warshall() { -----//7bc
- for (int k = 0; k < n; ++k) -----//7bd
-- for (int i = 0; i < n; ++i) -----//7be
--- for (int j = 0; j < n; ++j) -----//7bf
---- if (mat[i][k] + mat[k][j] < mat[i][j]) -----//7c0
----- mat[i][j] = mat[i][k] + mat[k][j]; -----//7c1
} -----//7c2
```

3.3. Strongly Connected Components.

3.3.1. Kosaraju.

```
struct kosaraju_graph { -----//74e
- int n; -----//74f
- int *vis; -----//750
- vi **adj; -----//751
- std::vector<vi> sccs; -----//752
- kosaraju_graph(int n) { -----//753
-- this->n = n; -----//754
-- vis = new int[n]; -----//755
-- adj = new vi*[2]; -----//756
-- for (int dir = 0; dir < 2; ++dir) -----//757
--- adj[dir] = new vi[n]; -----//758
- } -----//759
- void add_edge(int u, int v) { -----//75a
-- adj[0][u].push_back(v); -----//75b
-- adj[1][v].push_back(u); -----//75c
```

```
- } -----//75d
- void dfs(int u, int p, int dir, vi &topo) { -----//75e
-- vis[u] = 1; -----//75f
-- for (int v : adj[dir][u]) -----//760
--- if (!vis[v] && v != p) -----//761
---- dfs(v, u, dir, topo); -----//762
-- topo.push_back(u); -----//763
- } -----//764
- void kosaraju() { -----//765
-- vi topo; -----//766
-- for (int u = 0; u < n; ++u) vis[u] = 0; -----//767
-- for (int u = 0; u < n; ++u) -----//768
--- if (!vis[u]) -----//769
---- dfs(u, -1, 0, topo); -----//76a
-- for (int u = 0; u < n; ++u) vis[u] = 0; -----//76b
-- for (int i = n-1; i >= 0; --i) { -----//76c
--- if (!vis[topo[i]]) { -----//76d
---- sccs.push_back({}); -----//76e
---- dfs(topo[i], -1, 1, sccs.back()); -----//76f
--- } -----//770
-- } -----//771
- } -----//772
}; -----//773
```

3.3.2. Tarjan's Offline Algorithm.

```
int n, id[N], low[N], st[N], in[N], TOP, ID; -----//774
int scc[N], SCC_SIZE; // 0 <= scc[u] < SCC_SIZE -----//775
vector<int> adj[N]; // 0-based adjlist -----//776
void dfs(int u) { -----//777
-- id[u] = low[u] = ID++; -----//778
-- st[TOP++] = u; in[u] = 1; -----//779
-- for (int v : adj[u]) { -----//77a
--- if (id[v] == -1) { -----//77b
---- dfs(v); -----//77c
---- low[u] = min(low[u], low[v]); -----//77d
--- } else if (in[v] == 1) -----//77e
---- low[u] = min(low[u], id[v]); -----//77f
-- } -----//780
-- if (id[u] == low[u]) { -----//781
--- int sid = SCC_SIZE++; -----//782
--- do { -----//783
---- int v = st[--TOP]; -----//784
---- in[v] = 0; scc[v] = sid; -----//785
--- } while (st[TOP] != u); -----//786
-- } -----//787
void tarjan() { // call tarjan() to load SCC -----//788
-- memset(id, -1, sizeof(int) * n); -----//789
-- SCC_SIZE = ID = TOP = 0; -----//78a
-- for (int i = 0; i < n; ++i) -----//78b
--- if (id[i] == -1) dfs(i); } -----//78c
```

3.4. Minimum Mean Weight Cycle. Run this for each strongly connected component

```
double min_mean_cycle(vector<vector<pair<int,double>>> adj){
- int n = size(adj); double mn = INFINITY; -----//5b2
- vector<vector<double>> > arr(n+1, vector<double>(n, mn)); //5b3
- arr[0][0] = 0; -----//5b4
```

```
- rep(k,l,n+1) rep(j,0,n) iter(it,adj[j]) -----//5b5
-- arr[k][it->first] = min(arr[k][it->first], -----//5b6
--- it->second + arr[k-1][j]); -----//5b7
- rep(k,0,n) { -----//5b8
-- double mx = -INFINITY; -----//5b9
-- rep(i,0,n) mx = max(mx, (arr[n][i]-arr[k][i])/(n-k)); //5ba
-- mn = min(mn, mx); } -----//5bb
- return mn; } -----//5bc
```

3.5. Cut Points and Bridges.

```
vii bridges; -----//4d4
vi adj[MAXN], disc, low, articulation_points; -----//4d5
int TIME; -----//4d6
void bridges_artics (int u, int p) { -----//4d7
- disc[u] = low[u] = TIME++; -----//4d8
- int children = 0; -----//4d9
- bool has_low_child = false; -----//4da
- for (int v : adj[u]) { -----//4db
-- if (disc[v] == -1) { -----//4dc
--- bridges_artics(v, u); -----//4dd
--- children++; -----//4de
--- if (disc[u] < low[v]) -----//4df
---- bridges.push_back({u, v}); -----//4e0
--- if (disc[u] <= low[v]) -----//4e1
---- has_low_child = true; -----//4e2
--- low[u] = min(low[u], low[v]); -----//4e3
-- } else if (v != p) -----//4e4
--- low[u] = min(low[u], disc[v]); -----//4e5
- } -----//4e6
- if ((p == -1 && children >= 2) || -----//4e7
-- (p != -1 && has_low_child)) -----//4e8
-- articulation_points.push_back(u); -----//4e9
} -----//4ea
```

3.6. Biconnected Components.

3.6.1. Bridge Tree.

3.6.2. Block-Cut Tree.

3.7. Minimum Spanning Tree.

3.7.1. Kruskal.

```
#include "graph_template_edgelist.cpp" -----//72c
#include "union_find.cpp" -----//72d
// insert inside graph; needs n, and edges -----//72e
void kruskal(viii &res) { -----//72f
- viii().swap(res); // or use res.clear(); -----//730
- std::priority_queue<iii, viii, std::greater<iii> > pq; -----//731
- for (auto &edge : edges) -----//732
-- pq.push(edge); -----//733
- union_find uf(n); -----//734
- while (!pq.empty()) { -----//735
-- auto node = pq.top(); pq.pop(); -----//736
-- int u = node.second.first; -----//737
-- int v = node.second.second; -----//738
-- if (uf.unite(u, v)) -----//739
--- res.push_back(node); -----//73a
```

```
- } -----//73b
} -----//73c

3.7.2. Prim.
#include "graph_template_adjlist.cpp" -----//73d
// insert inside graph; needs n, vis[], and adj[] -----//73e
void prim(viii &res, int s=0) { -----//73f
- viii().swap(res); // or use res.clear(); -----//740
- std::priority_queue<ii, vii, std::greater<ii> > pq; -----//741
- pq.push({0, s}); -----//742
- while (!pq.empty()) { -----//743
-- int u = pq.top().second; pq.pop(); -----//744
-- vis[u] = true; -----//745
-- for (auto &[v, w] : adj[u]) { -----//746
---- if (v == u) continue; -----//747
---- if (vis[v]) continue; -----//748
---- res.push_back({w, {u, v}}); -----//749
---- pq.push({w, v}); -----//74a
-- } -----//74b
- } -----//74c
} -----//74d

3.8. Euler Path/Cycle.
3.8.1. Euler Path/Cycle in a Directed Graph.
#define MAXV 1000 -----//514
#define MAXE 5000 -----//515
vi adj[MAXV]; -----//516
int n, m, indeg[MAXV], outdeg[MAXV], res[MAXE + 1]; -----//517
ii start_end() { -----//518
- int start = -1, end = -1, any = 0, c = 0; -----//519
- rep(i,0,n) { -----//51a
-- if (outdeg[i] > 0) any = i; -----//51b
-- if (indeg[i] + 1 == outdeg[i]) start = i, c++; -----//51c
-- else if (indeg[i] == outdeg[i] + 1) end = i, c++; -----//51d
-- else if (indeg[i] != outdeg[i]) return ii(-1,-1); } -----//51e
- if ((start == -1) != (end == -1) || (c != 2 && c != 0)) -----//51f
-- return ii(-1,-1); -----//520
- if (start == -1) start = end = any; -----//521
- return ii(start, end); } -----//522
bool euler_path() { -----//523
- ii se = start_end(); -----//524
- int cur = se.first, at = m + 1; -----//525
- if (cur == -1) return false; -----//526
- stack<int> s; -----//527
- while (true) { -----//528
-- if (outdeg[cur] == 0) { -----//529
---- res[--at] = cur; -----//52a
---- if (s.empty()) break; -----//52b
---- cur = s.top(); s.pop(); -----//52c
-- } else s.push(cur), cur = adj[cur][--outdeg[cur]]; } -----//52d
- return at == 0; } -----//52e

3.8.2. (. Euler Path/Cycle in an Undirected Graph)
multiset<int> adj[1010]; -----//52f
list<int> L; -----//530
list<int>::iterator euler(int at, int to, -----//531
-- list<int>::iterator it) { -----//532
- if (at == to) return it; -----//533
- L.insert(it, at), --it; -----//534
- while (!adj[at].empty()) { -----//535
-- int nxt = *adj[at].begin(); -----//536
-- adj[at].erase(adj[at].find(nxt)); -----//537
-- adj[nxt].erase(adj[nxt].find(at)); -----//538
-- if (to == -1) { -----//539
--- it = euler(nxt, at, it); -----//53a
--- L.insert(it, at); -----//53b
--- --it; -----//53c
-- } else { -----//53d
--- it = euler(nxt, to, it); -----//53e
--- to = -1; } } -----//53f
- return it; } -----//540
// euler(0,-1,L.begin()) -----//541

3.9. Bipartite Matching.
3.9.1. Alternating Paths Algorithm.
vi* adj; -----//55e
bool* done; -----//55f
int* owner; -----//55f0
int alternating_path(int left) { -----//55f1
- if (done[left]) return 0; -----//55f2
- done[left] = true; -----//55f3
- rep(i,0,size(adj[left])) { -----//55f4
-- int right = adj[left][i]; -----//55f5
-- if (owner[right] == -1 || -----//55f6
--- alternating_path(owner[right])) { -----//55f7
--- owner[right] = left; return 1; } } -----//55f8
- return 0; } -----//55f9

3.9.2. Hopcroft-Karp Algorithm.
#define MAXN 5000 -----//609
int dist[MAXN+1], q[MAXN+1]; -----//60a
#define dist(v) dist[v == -1 ? MAXN : v] -----//60b
struct bipartite_graph { -----//60c
- int N, M, *L, *R; vi *adj; -----//60d
- bipartite_graph(int _N, int _M) : N(_N), M(_M), -----//60e
- L(new int[N]), R(new int[M]), adj(new vi[N]) {} -----//60f
- ~bipartite_graph() { delete[] adj; delete[] L; delete[] R; }
- bool bfs() { -----//611
-- int l = 0, r = 0; -----//612
-- rep(v,0,N) if(L[v] == -1) dist(v) = 0, q[r++] = v; -----//613
-- else dist(v) = INF; -----//614
-- dist(-1) = INF; -----//615
-- while(l < r) { -----//616
--- int v = q[l++]; -----//617
--- if(dist(v) < dist(-1)) { -----//618
--- iter(u, adj[v]) if(dist(R[*u]) == INF) -----//619
--- dist(R[*u]) = dist(v) + 1, q[r++] = R[*u]; } } -----//61a
-- return dist(-1) != INF; } -----//61b
- bool dfs(int v) { -----//61c
-- if(v != -1) { -----//61d
--- iter(u, adj[v]) -----//61e
--- if(dist(R[*u]) == dist(v) + 1) -----//61f
--- if(dfs(R[*u])) { -----//620
R[*u] = v, L[v] = *u; -----//621
return true; } -----//622
dist(v) = INF; -----//623
return false; } -----//624
return true; } -----//625
void add_edge(int i, int j) { adj[i].push_back(j); } -----//626
int maximum_matching() { -----//627
- int matching = 0; -----//628
- memset(L, -1, sizeof(int) * N); -----//629
- memset(R, -1, sizeof(int) * M); -----//62a
- while(bfs()) rep(i,0,N) -----//62b
-- matching += L[i] == -1 && dfs(i); -----//62c
- return matching; } } -----//62d

3.9.3. Minimum Vertex Cover in Bipartite Graphs.
#include "hopcroft_karp.cpp" -----//5fa
vector<bool> alt; -----//5fb
void dfs(bipartite_graph &g, int at) { -----//5fc
- alt[at] = true; -----//5fd
- iter(it,g.adj[at]) { -----//5fe
-- alt[*it + g.N] = true; -----//5ff
-- if (g.R[*it] != -1 && !alt[g.R[*it]]) -----//600
--- dfs(g, g.R[*it]); } } -----//601
vi mvc_bipartite(bipartite_graph &g) { -----//602
- vi res; g.maximum_matching(); -----//603
- alt.assign(g.N + g.M,false); -----//604
- rep(i,0,g.N) if (g.L[i] == -1) dfs(g, i); -----//605
- rep(i,0,g.N) if (!alt[i]) res.push_back(i); -----//606
- rep(i,0,g.M) if (alt[g.N + i]) res.push_back(g.N + i); -----//607
- return res; } -----//608

3.10. Maximum Flow.
3.10.1. Edmonds-Karp.
struct flow_network { -----//6b0
- int n, s, t, *par, **c, **f; -----//6b1
- vi *adj; -----//6b2
- flow_network(int n, int s, int t) : n(n), s(s), t(t) { -----//6b3
-- adj = new std::vector<int>[n]; -----//6b4
-- par = new int[n]; -----//6b5
-- c = new int*[n]; -----//6b6
-- f = new int*[n]; -----//6b7
-- for (int i = 0; i < n; ++i) { -----//6b8
--- c[i] = new int[n]; -----//6b9
--- f[i] = new int[n]; -----//6ba
--- for (int j = 0; j < n; ++j) -----//6bb
--- c[i][j] = f[i][j] = 0; -----//6bc
-- } -----//6bd
- } -----//6be
- void add_edge(int u, int v, int w) { -----//6bf
-- adj[u].push_back(v); -----//6c0
-- adj[v].push_back(u); -----//6c1
-- c[u][v] += w; -----//6c2
- } -----//6c3
- int res(int i, int j) { return c[i][j] - f[i][j]; } -----//6c4
- bool bfs() { -----//6c5
-- std::queue<int> q; -----//6c6
```

```
-- q.push(this->s); -----//6c7 - } -----//677 - } -----//6ae
-- while (!q.empty()) { -----//6c8 - int res(int i, int j) { return c[i][j] - f[i][j]; } -----//678 }; -----//6af
-- int u = q.front(); q.pop(); -----//6c9 - void reset(int *ar, int val) { -----//679
-- for (int v : adj[u]) { -----//6ca - for (int i = 0; i < n; ++i) -----//67a
-- if (res(u, v) > 0 and par[v] == -1) { -----//6cb - ar[i] = val; -----//67b
-- par[v] = u; -----//6cc - } -----//67c
-- if (v == this->t) -----//6cd - bool make_level_graph() { -----//67d
-- return true; -----//6ce - reset(dist, -1); -----//67e
-- q.push(v); -----//6cf - std::queue<int> q; -----//67f
-- } -----//6d0 - q.push(s); -----//680
-- } -----//6d1 - dist[s] = 0; -----//681
-- } -----//6d2 - while (!q.empty()) { -----//682
-- return false; -----//6d3 - int u = q.front(); q.pop(); -----//683
-- } -----//6d4 - for (int v : adj[u]) { -----//684
-- bool aug_path() { -----//6d5 - if (res(u, v) > 0 and dist[v] == -1) { -----//685
-- for (int u = 0; u < n; ++u) -----//6d6 - dist[v] = dist[u] + 1; -----//686
-- par[u] = -1; -----//6d7 - q.push(v); -----//687
-- par[s] = s; -----//6d8 - } -----//688
-- return bfs(); -----//6d9 - } -----//689
-- } -----//6da - } -----//68a
-- int calc_max_flow() { -----//6db - return dist[t] != -1; -----//68b
-- int ans = 0; -----//6dc - } -----//68c
-- while (aug_path()) { -----//6dd - bool next(int u, int v) { -----//68d
-- int flow = INF; -----//6de - return dist[v] == dist[u] + 1; -----//68e
-- for (int u = t; u != s; u = par[u]) -----//6df - } -----//68f
-- flow = std::min(flow, res(par[u], u)); -----//6e0 - bool dfs(int u) { -----//690
-- for (int u = t; u != s; u = par[u]) -----//6e1 - if (u == t) return true; -----//691
-- f[par[u]][u] += flow, f[u][par[u]] -= flow; -----//6e2 - for (int &i = adj_ptr[u]; i < adj[u].size(); ++i) { -----//692
-- ans += flow; -----//6e3 - int v = adj[u][i]; -----//693
-- } -----//6e4 - if (next(u, v) and res(u, v) > 0 and dfs(v)) { -----//694
-- return ans; -----//6e5 - par[v] = u; -----//695
-- } -----//6e6 - return true; -----//696
-- }; -----//6e7 - } -----//697
-- } -----//698
-- dist[u] = -1; -----//699
-- return false; -----//69a
-- } -----//69b
-- bool aug_path() { -----//69c
-- reset(par, -1); -----//69d
-- par[s] = s; -----//69e
-- return dfs(s); } -----//69f
-- int calc_max_flow() { -----//6a0
-- int ans = 0; -----//6a1
-- while (make_level_graph()) { -----//6a2
-- reset(adj_ptr, 0); -----//6a3
-- while (aug_path()) { -----//6a4
-- int flow = INF; -----//6a5
-- for (int u = t; u != s; u = par[u]) -----//6a6
-- flow = std::min(flow, res(par[u], u)); -----//6a7
-- for (int u = t; u != s; u = par[u]) -----//6a8
-- f[par[u]][u] += flow, f[u][par[u]] -= flow; -----//6a9
-- ans += flow; -----//6aa
-- } -----//6ab
-- } -----//6ac
-- return ans; -----//6ad

3.10.2. Dinic.
struct flow_network { -----//662
- int n, s, t, *adj_ptr, *dist, *par, **c, **f; -----//663
- vi *adj; -----//664
- flow_network(int n, int s, int t) : n(n), s(s), t(t) { -----//665
- adj = new std::vector<int>[n]; -----//666
- adj_ptr = new int[n]; -----//667
- dist = new int[n]; -----//668
- par = new int[n]; -----//669
- c = new int*[n]; -----//66a
- f = new int*[n]; -----//66b
- for (int i = 0; i < n; ++i) { -----//66c
- c[i] = new int[n]; -----//66d
- f[i] = new int[n]; -----//66e
- for (int j = 0; j < n; ++j) -----//66f
- c[i][j] = f[i][j] = 0; -----//670
- } -----//671
- } -----//672
- void add_edge(int u, int v, int w) { -----//673
- adj[u].push_back(v); -----//674
- adj[v].push_back(u); -----//675
- c[u][v] += w; -----//676
}
```

3.11. All-pairs Maximum Flow.

3.11.1. Gomory-Hu.

```
#define MAXV 2000 -----//6e8
int q[MAXV], d[MAXV]; -----//6e9
struct flow_network { -----//6ea
- struct edge { int v, nxt, cap; -----//6eb
- edge(int _v, int _cap, int _nxt) -----//6ec
- : v(_v), nxt(_nxt), cap(_cap) { } }; -----//6ed
- int n, *head, *curh; vector<edge> e, e_store; -----//6ee
- flow_network(int _n) : n(_n) { -----//6ef
- curh = new int[n]; -----//6f0
- memset(head = new int[n], -1, n*sizeof(int)); } -----//6f1
- void reset() { e = e_store; } -----//6f2
- void add_edge(int u, int v, int uv, int vu=0) { -----//6f3
- e.push_back(edge(v,uv,head[u])); head[u]=(int)size(e)-1;
- e.push_back(edge(u,vu,head[v])); head[v]=(int)size(e)-1;
- int augment(int v, int t, int f) { -----//6f6
- if (v == t) return f; -----//6f7
- for (int &i = curh[v], ret; i != -1; i = e[i].nxt) -----//6f8
- if (e[i].cap > 0 && d[e[i].v] + 1 == d[v]) -----//6f9
- if ((ret = augment(e[i].v, t, min(f, e[i].cap))) > 0)
- return (e[i].cap -= ret, e[i+1].cap += ret, ret);
- return 0; } -----//6fc
- int max_flow(int s, int t, bool res=true) { -----//6fd
- e_store = e; -----//6fe
- int l, r, f = 0, x; -----//6ff
- while (true) { -----//700
- memset(d, -1, n*sizeof(int)); -----//701
- l = r = 0, d[q[r++]] = t = 0; -----//702
- while (l < r) -----//703
- for (int v = q[l++], i = head[v]; i != -1; i=e[i].nxt)
- if (e[i+1].cap > 0 && d[e[i].v] == -1) -----//705
- d[q[r++]] = e[i].v = d[v]+1; -----//706
- if (d[s] == -1) break; -----//707
- memcpy(curh, head, n * sizeof(int)); -----//708
- while ((x = augment(s, t, INF)) != 0) f += x; } -----//709
- if (res) reset(); -----//70a
- return f; } }; -----//70b
bool same[MAXV]; -----//70c
pair<vii, vvii> construct_gh_tree(flow_network &g) { -----//70d
- int n = g.n, v; -----//70e
- vii par(n, ii(0, 0)); vvii cap(n, vi(n, -1)); -----//70f
- rep(s,1,n) { -----//710
- int l = 0, r = 0; -----//711
- par[s].second = g.max_flow(s, par[s].first, false); -----//712
- memset(d, 0, n * sizeof(int)); -----//713
- memset(same, 0, n * sizeof(bool)); -----//714
- d[q[r++]] = s = 1; -----//715
- while (l < r) { -----//716
- same[v = q[l++]] = true; -----//717
- for (int i = g.head[v]; i != -1; i = g.e[i].nxt) -----//718
- if (g.e[i].cap > 0 && d[g.e[i].v] == 0) -----//719
- d[q[r++]] = g.e[i].v = 1; } -----//71a
}
```



```
--- rep(i,s+1,n) -----//71b
---- if (par[i].first == par[s].first && same[i]) -----//71c
----- par[i].first = s; -----//71d
--- g.reset(); } -----//71e
- rep(i,0,n) { -----//71f
--- int mn = INF, cur = i; -----//720
--- while (true) { -----//721
----- cap[cur][i] = mn; -----//722
----- if (cur == 0) break; -----//723
----- mn = min(mn, par[cur].second), cur = par[cur].first; } }
- return make_pair(par, cap); } -----//725
int compute_max_flow(int s, int t, const pair<vii, vvi> &gh) {
- int cur = INF, at = s; -----//727
- while (gh.second[at][t] == -1) -----//728
-- cur = min(cur, gh.first[at].second), -----//729
-- at = gh.first[at].first; -----//72a
- return min(cur, gh.second[at][t]); } -----//72b
```

3.12. **Minimum Arborescence.** Given a weighted directed graph, finds a subset of edges of minimum total weight so that there is a unique path from the root r to each vertex. Returns a vector of size n , where the i th element is the edge for the i th vertex. The answer for the root is undefined!

```
#include "../data-structures/union_find.cpp" -----//45b
struct arborescence { -----//45c
- int n; union_find uf; -----//45d
- vector<vector<pair<ii,int> > > adj; -----//45e
- arborescence(int _n) : n(_n), uf(n), adj(n) { } -----//45f
- void add_edge(int a, int b, int c) { -----//460
-- adj[b].push_back(make_pair(ii(a,b),c)); } -----//461
- vii find_min(int r) { -----//462
-- vi vis(n,-1), mn(n,INF); vii par(n); -----//463
-- rep(i,0,n) { -----//464
----- if (uf.find(i) != i) continue; -----//465
----- int at = i; -----//466
----- while (at != r && vis[at] == -1) { -----//467
----- vis[at] = i; -----//468
----- iter(it,adj[at]) if (it->second < mn[at] && -----//469
----- uf.find(it->first.first) != at) -----//46a
----- mn[at] = it->second, par[at] = it->first; -----//46b
----- if (par[at] == ii(0,0)) return vii(); -----//46c
----- at = uf.find(par[at].first); } -----//46d
----- if (at == r || vis[at] != i) continue; -----//46e
----- union_find tmp = uf; vi seq; -----//46f
----- do { seq.push_back(at); at = uf.find(par[at].first); } -----//470
----- } while (at != seq.front()); -----//471
----- iter(it,seq) uf.unite(*it,seq[0]); -----//472
----- int c = uf.find(seq[0]); -----//473
----- vector<pair<ii,int> > nw; -----//474
----- iter(it,seq) iter(jt,adj[*it]) -----//475
----- nw.push_back(make_pair(jt->first, -----//476
----- jt->second - mn[*it])); -----//477
----- adj[c] = nw; -----//478
----- vii rest = find_min(r); -----//479
----- if (size(rest) == 0) return rest; -----//47a
----- ii use = rest[c]; -----//47b
```

```
rest[at = tmp.find(use.second)] = use; -----//47c
iter(it,seq) if (*it != at) -----//47d
rest[*it] = par[*it]; -----//47e
return rest; } -----//47f
return par; } };
```

3.13. **Blossom algorithm.** Finds a maximum matching in an arbitrary graph in $O(|V|^4)$ time. Be vary of loop edges.

```
#define MAXV 300 -----//481
bool marked[MAXV], emarked[MAXV][MAXV]; -----//482
int S[MAXV]; -----//483
vi find_augmenting_path(const vector<vi> &adj, const vi &m){
- int n = size(adj), s = 0; -----//485
- vi par(n,-1), height(n), root(n,-1), q, a, b; -----//486
- memset(marked,0,sizeof(marked)); -----//487
- memset(emarked,0,sizeof(emarked)); -----//488
- rep(i,0,n) if (m[i] >= 0) emarked[i][m[i]] = true; -----//489
----- else root[i] = i, S[s++] = i; -----//48a
- while (s) { -----//48b
-- int v = S[--s]; -----//48c
-- iter(wt,adj[v]) { -----//48d
--- int w = *wt; -----//48e
--- if (emarked[v][w]) continue; -----//48f
--- if (root[w] == -1) { -----//490
---- int x = S[s++] = m[w]; -----//491
---- par[w]=v, root[w]=root[v], height[w]=height[v]+1; -----//492
---- par[x]=w, root[x]=root[w], height[x]=height[w]+1; -----//493
---- } else if (height[w] % 2 == 0) { -----//494
---- if (root[v] != root[w]) { -----//495
----- while (v != -1) q.push_back(v), v = par[v]; -----//496
----- reverse(q.begin(), q.end()); -----//497
----- while (w != -1) q.push_back(w), w = par[w]; -----//498
----- return q; -----//499
---- } else { -----//49a
---- int c = v; -----//49b
---- while (c != -1) a.push_back(c), c = par[c]; -----//49c
---- c = w; -----//49d
---- while (c != -1) b.push_back(c), c = par[c]; -----//49e
---- while (!a.empty()&&!b.empty()&&a.back()==b.back())
-- c = a.back(), a.pop_back(), b.pop_back(); -----//4a0
-- memset(marked,0,sizeof(marked)); -----//4a1
-- fill(par.begin(), par.end(), 0); -----//4a2
-- iter(it,a) par[*it] = 1; iter(it,b) par[*it] = 1;
-- par[c] = s = 1; -----//4a4
-- rep(i,0,n) root[par[i] = par[i] ? 0 : s++] = i; -----//4a5
-- vector<vi> adj2(s); -----//4a6
-- rep(i,0,n) iter(it,adj[i]) { -----//4a7
--- if (par[*it] == 0) continue; -----//4a8
--- if (par[i] == 0) { -----//4a9
---- if (!marked[par[*it]]) { -----//4aa
----- adj2[par[i]].push_back(par[*it]); -----//4ab
----- adj2[par[*it]].push_back(par[i]); -----//4ac
----- marked[par[*it]] = true; } -----//4ad
---- } else adj2[par[i]].push_back(par[*it]); } -----//4ae
-- vi m2(s, -1); -----//4af
-- if (m[c] != -1) m2[m2[par[m[c]]] = 0] = par[m[c]];
```

```
rep(i,0,n) if(par[i]!=0&&m[i]!=-1&&par[m[i]]!=0)//4b1
m2[par[i]] = par[m[i]]; -----//4b2
vi p = find_augmenting_path(adj2, m2); -----//4b3
int t = 0; -----//4b4
while (t < size(p) && p[t]) t++; -----//4b5
if (t == size(p)) { -----//4b6
rep(i,0,size(p)) p[i] = root[p[i]]; -----//4b7
return p; } -----//4b8
if (!p[i] || (m[c] != -1 && p[t+1] != par[m[c]]))
reverse(p.begin(), p.end()), t=(int)size(p)-t-1;
rep(i,0,t) q.push_back(root[p[i]]); -----//4bb
iter(it,adj[root[p[t-1]]) { -----//4bc
if (par[*it] != (s = 0)) continue; -----//4bd
a.push_back(c), reverse(a.begin(), a.end()); -----//4be
iter(jt,b) a.push_back(*jt); -----//4bf
while (a[s] != *it) s++; -----//4c0
if((height[*it]&1)^(s<(int)size(a)-(int)size(b)))
reverse(a.begin(),a.end()), s=(int)size(a)-s-1;
while(a[s]!=c)q.push_back(a[s]),s=(s+1)%size(a);
q.push_back(c); -----//4c4
rep(i,t+1,size(p)) q.push_back(root[p[i]]); -----//4c5
return q; } } } -----//4c6
emarked[v][w] = emarked[w][v] = true; } -----//4c7
marked[v] = true; } return q; } -----//4c8
vii max_matching(const vector<vi> &adj) { -----//4c9
- vi m(size(adj), -1), ap; vii res, es; -----//4ca
- rep(i,0,size(adj)) iter(it,adj[i]) es.emplace_back(i,*it);
- random_shuffle(es.begin(), es.end()); -----//4cb
- iter(it,es) if (m[it->first] == -1 && m[it->second] == -1)
-- m[it->first] = it->second, m[it->second] = it->first; -----//4ce
-- do { ap = find_augmenting_path(adj, m); -----//4cf
----- rep(i,0,size(ap)) m[m[ap[i^1]] = ap[i]] = ap[i^1]; -----//4d0
-- } while (!ap.empty()); -----//4d1
- rep(i,0,size(m)) if (i < m[i]) res.emplace_back(i, m[i]);
- return res; } -----//4d3
```

3.14. **Maximum Density Subgraph.** Given (weighted) undirected graph G . Binary search density. If g is current density, construct flow network: $(S, u, m), (u, T, m + 2g - d_u), (u, v, 1)$, where m is a large constant (larger than sum of edge weights). Run floating-point max-flow. If minimum cut has empty S -component, then maximum density is smaller than g , otherwise it's larger. Distance between valid densities is at least $1/(n(n-1))$. Edge case when density is 0. This also works for weighted graphs by replacing d_u by the weighted degree, and doing more iterations (if weights are not integers).

3.15. **Maximum-Weight Closure.** Given a vertex-weighted directed graph G . Turn the graph into a flow network, adding weight ∞ to each edge. Add vertices S, T . For each vertex v of weight w , add edge (S, v, w) if $w \geq 0$, or edge $(v, T, -w)$ if $w < 0$. Sum of positive weights minus minimum $S - T$ cut is the answer. Vertices reachable from S are in the closure. The maximum-weight closure is the same as the complement of the minimum-weight closure on the graph with edges reversed.

3.16. **Maximum Weighted Independent Set in a Bipartite Graph.** This is the same as the minimum weighted vertex cover. Solve this by constructing a flow network with edges $(S, u, w(u))$ for $u \in L$,

$(v, T, w(v))$ for $v \in R$ and (u, v, ∞) for $(u, v) \in E$. The minimum S, T -cut is the answer. Vertices adjacent to a cut edge are in the vertex cover.

3.17. Synchronizing word problem. A DFA has a synchronizing word (an input sequence that moves all states to the same state) iff. each pair of states has a synchronizing word. That can be checked using reverse DFS over pairs of states. Finding the shortest synchronizing word is NP-complete.

3.18. Max flow with lower bounds on edges. Change edge $(u, v, l \leq f \leq c)$ to $(u, v, f \leq c - l)$. Add edge (t, s, ∞) . Create super-nodes S, T . Let $M(u) = \sum_v l(v, u) - \sum_v l(u, v)$. If $M(u) < 0$, add edge $(u, T, -M(u))$, else add edge $(S, u, M(u))$. Max flow from S to T . If all edges from S are saturated, then we have a feasible flow. Continue running max flow from s to t in original graph.

3.19. Tutte matrix for general matching. Create an $n \times n$ matrix A . For each edge (i, j) , $i < j$, let $A_{ij} = x_{ij}$ and $A_{ji} = -x_{ij}$. All other entries are 0. The determinant of A is zero iff. the graph has a perfect matching. A randomized algorithm uses the Schwartz–Zippel lemma to check if it is zero.

3.20. Heavy Light Decomposition.

```
#include "segment_tree.cpp" -----//568
struct heavy_light_tree { -----//569
-   int n; -----//56a
-   std::vector<int> *adj; -----//56b
-   segtree *segment_tree; -----//56c
-   int *par, *heavy, *dep, *path_root, *pos; -----//56d
-   heavy_light_tree(int n) { -----//56e
-   this->n = n; -----//56f
-   this->adj = new std::vector<int>[n]; -----//570
-   segment_tree = new segtree(0, n-1); -----//571
-   par = new int[n]; -----//572
-   heavy = new int[n]; -----//573
-   dep = new int[n]; -----//574
-   path_root = new int[n]; -----//575
-   pos = new int[n]; -----//576
-   } -----//577
-   void add_edge(int u, int v) { -----//578
-   adj[u].push_back(v); -----//579
-   adj[v].push_back(u); -----//57a
-   } -----//57b
-   void build(int root) { -----//57c
-   for (int u = 0; u < n; ++u) -----//57d
-   heavy[u] = -1; -----//57e
-   par[root] = root; -----//57f
-   dep[root] = 0; -----//580
-   dfs(root); -----//581
-   for (int u = 0, p = 0; u < n; ++u) { -----//582
-   if (par[u] == -1 or heavy[par[u]] != u) { -----//583
-   for (int v = u; v != -1; v = heavy[v]) { -----//584
-   path_root[v] = u; -----//585
-   pos[v] = p++; -----//586
-   } -----//587
-   } -----//588
-   } -----//589
-   } -----//58a
```

```
-   int dfs(int u) { -----//58b
-   int sz = 1; -----//58c
-   int max_subtree_sz = 0; -----//58d
-   for (int v : adj[u]) { -----//58e
-   if (v != par[u]) { -----//58f
-   par[v] = u; -----//590
-   dep[v] = dep[u] + 1; -----//591
-   int subtree_sz = dfs(v); -----//592
-   if (max_subtree_sz < subtree_sz) { -----//593
-   max_subtree_sz = subtree_sz; -----//594
-   heavy[u] = v; -----//595
-   } -----//596
-   sz += subtree_sz; -----//597
-   } -----//598
-   } -----//599
-   return sz; -----//59a
-   } -----//59b
-   int query(int u, int v) { -----//59c
-   int res = 0; -----//59d
-   while (path_root[u] != path_root[v]) { -----//59e
-   if (dep[path_root[u]] > dep[path_root[v]]) -----//59f
-   std::swap(u, v); -----//5a0
-   res += segment_tree->sum(pos[path_root[v]], pos[v]); //5a1
-   v = par[path_root[v]]; -----//5a2
-   } -----//5a3
-   res += segment_tree->sum(pos[u], pos[v]); -----//5a4
-   return res; -----//5a5
-   } -----//5a6
-   void update(int u, int v, int c) { -----//5a7
-   for (; path_root[u] != path_root[v]; -----//5a8
-   v = par[path_root[v]]) { -----//5a9
-   if (dep[path_root[u]] > dep[path_root[v]]) -----//5aa
-   std::swap(u, v); -----//5ab
-   segment_tree->increase(pos[path_root[v]], pos[v], c); -----//5ad
-   } -----//5ae
-   segment_tree->increase(pos[u], pos[v], c); -----//5af
-   } -----//5b0
}; -----//5b0
```

3.21. Centroid Decomposition.

```
#define MAXV 100100 -----//4eb
#define LGMAXV 20 -----//4ec
int jmp[MAXV][LGMAXV], -----//4ed
-   path[MAXV][LGMAXV], -----//4ee
-   sz[MAXV], seph[MAXV], -----//4ef
-   shortest[MAXV]; -----//4f0
struct centroid_decomposition { -----//4f1
-   int n; vvi adj; -----//4f2
-   centroid_decomposition(int _n) : n(_n), adj(n) { } -----//4f3
-   void add_edge(int a, int b) { -----//4f4
-   adj[a].push_back(b); adj[b].push_back(a); } -----//4f5
-   int dfs(int u, int p) { -----//4f6
-   sz[u] = 1; -----//4f7
-   rep(i,0,size(adj[u])) -----//4f8
-   if (adj[u][i] != p) sz[u] += dfs(adj[u][i], u); -----//4f9
-   return sz[u]; } -----//4fa
```

```
-   void makepaths(int sep, int u, int p, int len) { -----//4fb
-   jmp[u][seph[sep]] = sep, path[u][seph[sep]] = len; ---//4fc
-   int bad = -1; -----//4fd
-   rep(i,0,size(adj[u])) { -----//4fe
-   if (adj[u][i] == p) bad = i; -----//4ff
-   else makepaths(sep, adj[u][i], u, len + 1); -----//500
-   } -----//501
-   if (p == sep) -----//502
-   swap(adj[u][bad], adj[u].back(), adj[u].pop_back()); } -----//504
-   void separate(int h=0, int u=0) { -----//50a
-   dfs(u, -1); int sep = u; -----//505
-   down: iter(nxt,adj[sep]) -----//506
-   if (sz[*nxt] < sz[sep] && sz[*nxt] > sz[u]/2) { ----//507
-   sep = *nxt; goto down; } -----//508
-   seph[sep] = h, makepaths(sep, sep, -1, 0); -----//509
-   rep(i,0,size(adj[sep])) separate(h+1, adj[sep][i]); } //50a
-   void paint(int u) { -----//50b
-   rep(h,0,seph[u]+1) -----//50c
-   shortest[jmp[u][h]] = min(shortest[jmp[u][h]], -----//50d
-   path[u][h]); } -----//50e
-   int closest(int u) { -----//50f
-   int mn = INF/2; -----//510
-   rep(h,0,seph[u]+1) -----//511
-   mn = min(mn, path[u][h] + shortest[jmp[u][h]]); ---//512
-   return mn; } }; -----//513
```

3.22. Least Common Ancestor.

3.22.1. Binary Lifting.

```
struct graph { -----//62e
-   int n; -----//62f
-   int logn; -----//630
-   std::vector<int> *adj; -----//631
-   int *dep; -----//632
-   int **par; -----//633
-   graph(int n, int logn=20) { -----//634
-   this->n = n; -----//635
-   this->logn = logn; -----//636
-   adj = new std::vector<int>[n]; -----//637
-   dep = new int[n]; -----//638
-   par = new int*[n]; -----//639
-   for (int i = 0; i < n; ++i) -----//63a
-   par[i] = new int[logn]; -----//63b
-   } -----//63c
-   void dfs(int u, int p, int d) { -----//63d
-   dep[u] = d; -----//63e
-   par[u][0] = p; -----//63f
-   for (int v : adj[u]) -----//640
-   if (v != p) -----//641
-   dfs(v, u, d+1); -----//642
-   } -----//643
-   int ascend(int u, int k) { -----//644
-   for (int i = 0; i < logn; ++i) -----//645
-   if (k & (1 << i)) -----//646
-   u = par[u][i]; -----//647
-   return u; -----//648
-   } -----//649
```

```
- int lca(int u, int v) { -----//64a
-- if (dep[u] > dep[v]) u = ascend(u, dep[u] - dep[v]); //64b
-- if (dep[v] > dep[u]) v = ascend(v, dep[v] - dep[u]); //64c
-- if (u == v) return u; -----//64d
-- for (int k = logn-1; k >= 0; --k) { -----//64e
---- if (par[u][k] != par[v][k]) { -----//64f
----- u = par[u][k]; -----//650
----- v = par[v][k]; -----//651
----- } -----//652
-- } -----//653
-- return par[u][0]; -----//654
- } -----//655
- bool is_anc(int u, int v) { -----//656
-- if (dep[u] < dep[v]) -----//657
----- std::swap(u, v); -----//658
-- return ascend(u, dep[u] - dep[v]) == v; -----//659
- } -----//65a
- void prep_lca(int root=0) { -----//65b
-- dfs(root, root, 0); -----//65c
-- for (int k = 1; k < logn; ++k) -----//65d
---- for (int u = 0; u < n; ++u) -----//65e
----- par[u][k] = par[par[u][k-1]][k-1]; -----//65f
- } -----//660
}; -----//661
```

3.23. **Counting Spanning Trees.** Kirchoff's Theorem: The number of spanning trees of any graph is the determinant of any cofactor of the Laplacian matrix in $O(n^3)$.

- (1) Let A be the adjacency matrix.
- (2) Let D be the degree matrix (matrix with vertex degrees on the diagonal).
- (3) Get $D - A$ and delete exactly one row and column. Any row and column will do. This will be the cofactor matrix.
- (4) Get the determinant of this cofactor matrix using Gauss-Jordan.
- (5) Spanning Trees = |cofactor($D - A$)|

3.24. **Erdős-Gallai Theorem.** A sequence of non-negative integers $d_1 \geq \dots \geq d_n$ can be represented as the degree sequence of finite simple graph on n vertices if and only if $d_1 + \dots + d_n$ is even and the following holds for $1 \leq k \leq n$:

$$\sum_{i=1}^n d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i, k)$$

3.25. **Tree Isomorphism.**

```
// REQUIREMENT: list of primes pr[], see prime sieve -----//5bd
typedef long long LL; -----//5be
int pre[N], q[N], path[N]; bool vis[N]; -----//5bf
// perform BFS and return the last node visited -----//5c0
int bfs(int u, vector<int> adj[]) { -----//5c1
-- memset(vis, 0, sizeof(vis)); -----//5c2
-- int head = 0, tail = 0; -----//5c3
-- q[tail++] = u; vis[u] = true; pre[u] = -1; -----//5c4
-- while (head != tail) { -----//5c5
----- u = q[head]; if (++head == N) head = 0; -----//5c6
----- for (int i = 0; i < adj[u].size(); ++i) { -----//5c7
----- int v = adj[u][i]; -----//5c8
```

```
----- if (!vis[v]) { -----//5c9
----- vis[v] = true; pre[v] = u; -----//5ca
----- q[tail++] = v; if (tail == N) tail = 0; -----//5cb
----- } -----//5cc
-- return u; -----//5cd
} // returns the list of tree centers -----//5ce
vector<int> tree_centers(int r, vector<int> adj[]) { -----//5cf
-- int size = 0; -----//5d0
-- for (int u=bfs(bfs(r, adj), adj); u!=-1; u=pre[u]) -----//5d1
---- path[size++] = u; -----//5d2
-- vector<int> med(1, path[size/2]); -----//5d3
-- if (size % 2 == 0) med.push_back(path[size/2-1]); -----//5d4
-- return med; -----//5d5
} // returns "unique hashcode" for tree with root u -----//5d6
LL rootcode(int u, vector<int> adj[], int p=-1, int d=15){//5d7
-- vector<LL> k; int nd = (d + 1) % primes; -----//5d8
-- for (int i = 0; i < adj[u].size(); ++i) -----//5d9
---- if (adj[u][i] != p) -----//5da
----- k.push_back(rootcode(adj[u][i], adj, u, nd)); //5db
-- sort(k.begin(), k.end()); -----//5dc
-- LL h = k.size() + 1; -----//5dd
-- for (int i = 0; i < k.size(); ++i) -----//5de
---- h = h * pr[d] + k[i]; -----//5df
-- return h; -----//5e0
} // returns "unique hashcode" for the whole tree -----//5e1
LL treecode(int root, vector<int> adj[]) { -----//5e2
-- vector<int> c = tree_centers(root, adj); -----//5e3
-- if (c.size()==1) -----//5e4
---- return (rootcode(c[0], adj) << 1) | 1; -----//5e5
-- return (rootcode(c[0],adj)*rootcode(c[1],adj))<<1; ---//5e6
} // checks if two trees are isomorphic -----//5e7
bool isomorphic(int r1, vector<int> adj1[], int r2, -----//5e8
----- vector<int> adj2[], bool rooted = false) { //5e9
-- if (rooted) -----//5ea
---- return rootcode(r1, adj1) == rootcode(r2, adj2); //5eb
-- return treecode(r1, adj1) == treecode(r2, adj2); ---//5ec
} -----//5ed
```

4. STRINGS

4.1. **Knuth-Morris-Pratt.** Count and find all matches of string f in string s in $O(n)$ time.

```
int par[N]; // parent table -----//8db
void buildKMP(string& f) { -----//8dc
-- par[0] = -1, par[1] = 0; -----//8dd
-- int i = 2, j = 0; -----//8de
-- while (i <= f.length()) { -----//8df
---- if (f[i-1] == f[j]) par[i++] = ++j; -----//8e0
---- else if (j > 0) j = par[j]; -----//8e1
---- else par[i++] = 0; } } -----//8e2
vector<int> KMP(string& s, string& f) { -----//8e3
-- buildKMP(f); // call once if f is the same -----//8e4
-- int i = 0, j = 0; vector<int> ans; -----//8e5
-- while (i + j < s.length()) { -----//8e6
---- if (s[i + j] == f[j]) { -----//8e7
----- if (++j == f.length()) { -----//8e8
----- ans.push_back(i); -----//8e9
```

```
----- i += j - par[j]; -----//8ea
----- if (j > 0) j = par[j]; -----//8eb
----- } -----//8ec
-- } else { -----//8ed
---- i += j - par[j]; -----//8ee
---- if (j > 0) j = par[j]; -----//8ef
---- } -----//8f0
-- } return ans; } -----//8f1
```

4.2. **Suffix Array.** Construct a sorted catalog of all substrings of s in $O(n \log n)$ time using counting sort.

```
// sa[i]: ith smallest substring at s[sa[i]:] -----//94c
// pos[i]: position of s[i:] in suffix array -----//94d
int sa[N], pos[N], va[N], c[N], gap, n; -----//94e
bool cmp(int i, int j) // reverse stable sort -----//94f
-- {return pos[i]!=pos[j] ? pos[i] < pos[j] : j < i;} ---//950
bool equal(int i, int j) -----//951
-- {return pos[i] == pos[j] && i + gap < n && -----//952
---- pos[i + gap / 2] == pos[j + gap / 2];} -----//953
void buildSA(string s) { -----//954
-- s += '$'; n = s.length(); -----//955
-- for (int i = 0; i < n; i++){sa[i]=i; pos[i]=s[i];} ---//956
-- sort (sa, sa + n, cmp); -----//957
-- for (gap = 1; gap < n * 2; gap <= 1) { -----//958
---- va[sa[0]] = 0; -----//959
---- for (int i = 1; i < n; i++) { -----//95a
----- int prev = sa[i - 1], next = sa[i]; -----//95b
----- va[next] = equal(prev, next) ? va[prev] : i; //95c
----- } -----//95d
---- for (int i = 0; i < n; ++i) -----//95e
----- { pos[i] = va[i]; va[i] = sa[i]; c[i] = i; } //95f
---- for (int i = 0; i < n; i++) { -----//960
----- int id = va[i] - gap; -----//961
----- if (id >= 0) sa[c[pos[id]]++] = id; -----//962
----- } } -----//963
```

4.3. **Longest Common Prefix.** Find the length of the longest common prefix for every substring in $O(n)$.

```
int lcp[N]; // lcp[i] = LCP(s[sa[i]:], s[sa[i+1]:]) -----//8f2
void buildLCP(string s) { // build suffix array first -----//8f3
-- for (int i = 0, k = 0; i < n; i++) { -----//8f4
---- if (pos[i] != n - 1) { -----//8f5
----- for(int j = sa[pos[i]+1]; s[i+k]==s[j+k];k++); //8f6
----- lcp[pos[i]] = k; if (k > 0) k--; -----//8f7
---- } else { lcp[pos[i]] = 0; } } -----//8f8
```

4.4. **Aho-Corasick Trie.** Find all multiple pattern matches in $O(n)$ time. This is KMP for multiple strings.

```
class Node { -----//89a
-- HashMap<Character, Node> next = new HashMap<>(); -----//89b
-- Node fail = null; -----//89c
-- long count = 0; -----//89d
-- public void add(String s) { // adds string to trie ---//89e
---- Node node = this; -----//89f
---- for (char c : s.toCharArray()) { -----//8a0
----- if (!node.contains(c)) -----//8a1
----- node.next.put(c, new Node()); -----//8a2
```

```
----- node = node.get(c); -----//8a3
----- } node.count++; -----//8a4
public void prepare() { -----//8a5
    // prepares fail links of Aho-Corasick Trie -----//8a6
    Node root = this; root.fail = null; -----//8a7
    Queue<Node> q = new ArrayDeque<Node>(); -----//8a8
    for (Node child : next.values()) // BFS -----//8a9
    { child.fail = root; q.offer(child); } -----//8aa
    while (!q.isEmpty()) { -----//8ab
        Node head = q.poll(); -----//8ac
        for (Character letter : head.next.keySet()) { //8ad
            // traverse upwards to get nearest fail link --//8ae
            Node p = head; -----//8af
            Node nextNode = head.get(letter); -----//8b0
            do { p = p.fail; } -----//8b1
            while(p != root && !p.contains(letter)); --//8b2
            if (p.contains(letter)) { // fail link found
                p = p.get(letter); -----//8b4
                nextNode.fail = p; -----//8b5
                nextNode.count += p.count; -----//8b6
            } else { nextNode.fail = root; } -----//8b7
            q.offer(nextNode); -----//8b8
        } } -----//8b9
    public BigInteger search(String s) { -----//8ba
        // counts the words added in trie present in s ---//8bb
        Node root = this, p = this; -----//8bc
        BigInteger ans = BigInteger.ZERO; -----//8bd
        for (char c : s.toCharArray()) { -----//8be
            while (p != root && !p.contains(c)) p = p.fail;
            if (p.contains(c)) { -----//8c0
                p = p.get(c); -----//8c1
                ans = ans.add(BigInteger.valueOf(p.count));
            } -----//8c3
        } return ans; } -----//8c4
    // helper methods -----//8c5
    private Node get(char c) { return next.get(c); } -----//8c6
    private boolean contains(char c) { -----//8c7
        return next.containsKey(c); -----//8c8
    } // Usage: Node trie = new Node(); -----//8c9
    // for (String s : dictionary) trie.add(s); -----//8ca
    // trie.prepare(); BigInteger m = trie.search(str); -----//8cb
```

4.5. **Palindromic Tree.** Find lengths and frequencies of all palindromic substrings of a string in $O(n)$ time.

Theorem: there can only be up to n unique palindromic substrings for any string.

```
int par[N*2+1], child[N*2+1][128]; -----//8f9
int len[N*2+1], node[N*2+1], cs[N*2+1], size; -----//8fa
long long cnt[N + 2]; // count can be very large -----//8fb
int newNode(int p = -1) { -----//8fc
    cnt[size] = 0; par[size] = p; -----//8fd
    len[size] = (p == -1 ? 0 : len[p] + 2); -----//8fe
    memset(child[size], -1, sizeof child[size]); -----//8ff
    return size++; -----//900
} -----//901
int get(int i, char c) { -----//902
```

```
    if (child[i][c] == -1) child[i][c] = newNode(i); -----//903
    return child[i][c]; -----//904
} -----//905
void manachers(char s[]) { -----//906
    int n = strlen(s), cn = n * 2 + 1; -----//907
    for (int i = 0; i < n; i++) -----//908
        {cs[i * 2] = -1; cs[i * 2 + 1] = s[i];} -----//909
    size = n * 2; -----//90a
    int odd = newNode(), even = newNode(); -----//90b
    int cen = 0, rad = 0, L = 0, R = 0; -----//90c
    size = 0; len[odd] = -1; -----//90d
    for (int i = 0; i < cn; i++) -----//90e
        node[i] = (i % 2 == 0 ? even : get(odd, cs[i])); --//90f
    for (int i = 1; i < cn; i++) { -----//910
        if (i > rad) { L = i - 1; R = i + 1; } -----//911
        else { -----//912
            int M = cen * 2 - i; // retrieve from mirror --//913
            node[i] = node[M]; -----//914
            if (len[node[M]] < rad - i) L = -1; -----//915
            else { -----//916
                R = rad + 1; L = i * 2 - R; -----//917
                while (len[node[i]] > rad - i) -----//918
                    node[i] = par[node[i]]; -----//919
            } -----//91a
        } // expand palindrome -----//91b
        while (L >= 0 && R < cn && cs[L] == cs[R]) { -----//91c
            if (cs[L] != -1) node[i] = get(node[i], cs[L]); --//91d
            L--, R++; -----//91e
        } -----//91f
        cnt[node[i]]++; -----//920
        if (i + len[node[i]] > rad) -----//921
            { rad = i + len[node[i]]; cen = i; } -----//922
        } -----//923
    for (int i = size - 1; i >= 0; --i) -----//924
        cnt[par[i]] += cnt[i]; // update parent count -----//925
} -----//926
int countUniquePalindromes(char s[]) -----//927
{manachers(s); return size;} -----//928
int countAllPalindromes(char s[]) { -----//929
    manachers(s); int total = 0; -----//92a
    for (int i = 0; i < size; i++) total += cnt[i]; -----//92b
    return total;} -----//92c
// longest palindrome substring of s -----//92d
string longestPalindrome(char s[]) { -----//92e
    manachers(s); -----//92f
    int n = strlen(s), cn = n * 2 + 1, mx = 0; -----//930
    for (int i = 1; i < cn; i++) -----//931
        if (len[node[mx]] < len[node[i]]) -----//932
            mx = i; -----//933
    int pos = (mx - len[node[mx]]) / 2; -----//934
    return string(s + pos, s + pos + len[node[mx]]); } --//935
4.6. Z Algorithm. Find the longest common prefix of all substrings of
s with itself in  $O(n)$  time.
int z[N]; // z[i] = lcp(s, s[i:]) -----//964
void computeZ(string s) { -----//965
```

```
    int n = s.length(), L = 0, R = 0; z[0] = n; -----//966
    for (int i = 1; i < n; i++) { -----//967
        if (i > R) { -----//968
            L = R = i; -----//969
            while (R < n && s[R - L] == s[R]) R++; -----//96a
            z[i] = R - L; R--; -----//96b
        } else { -----//96c
            int k = i - L; -----//96d
            if (z[k] < R - i + 1) z[i] = z[k]; -----//96e
            else { -----//96f
                L = i; -----//970
                while (R < n && s[R - L] == s[R]) R++; --//971
                z[i] = R - L; R--; -----//972
            } } } -----//973
```

4.7. **Booth's Minimum String Rotation.** Booth's Algo: Find the index of the lexicographically least string rotation in $O(n)$ time.

```
int f[N * 2]; -----//8cc
int booth(string S) { -----//8cd
    S.append(S); // concatenate itself -----//8ce
    int n = S.length(), i, j, k = 0; -----//8cf
    memset(f, -1, sizeof(int) * n); -----//8d0
    for (j = 1; j < n; j++) { -----//8d1
        i = f[j-k-1]; -----//8d2
        while (i != -1 && S[j] != S[k + i + 1]) { -----//8d3
            if (S[j] < S[k + i + 1]) k = j - i - 1; -----//8d4
            i = f[i]; -----//8d5
        } if (i == -1 && S[j] != S[k + i + 1]) { -----//8d6
            if (S[j] < S[k + i + 1]) k = j; -----//8d7
            f[j - k] = -1; -----//8d8
        } else f[j - k] = i + 1; -----//8d9
    } return k; } -----//8da
```

4.8. **Hashing.**

4.8.1. *Polynomial Hashing.*

```
int MAXN = 1e5+1, MOD = 1e9+7; -----//936
struct hasher { -----//937
    int n; -----//938
    std::vector<ll> *p_pow; -----//939
    std::vector<ll> *h_ans; -----//93a
    hash(vi &s, vi primes) { -----//93b
        n = primes.size(); -----//93c
        p_pow = new std::vector<ll>[n]; -----//93d
        h_ans = new std::vector<ll>[n]; -----//93e
        for (int i = 0; i < n; ++i) { -----//93f
            p_pow[i] = std::vector<ll>(MAXN); -----//940
            p_pow[i][0] = 1; -----//941
            for (int j = 0; j+1 < MAXN; ++j) -----//942
                p_pow[i][j+1] = (p_pow[i][j] * primes[i]) % MOD; --//943
            h_ans[i] = std::vector<ll>(MAXN); -----//944
            h_ans[i][0] = 0; -----//945
            for (int j = 0; j < s.size(); ++j) -----//946
                h_ans[i][j+1] = (h_ans[i][j] + -----//947
                    s[j] * p_pow[i][j]) % MOD; -----//948
        } -----//949
```



```
- } -----//94a
}; -----//94b
```

5. DYNAMIC PROGRAMMING

5.1. Longest Common Subsequence.

5.2. Longest Increasing Subsequence.

5.3. Traveling Salesman.

6. NUMBER THEORY

6.1. Eratosthenes Prime Sieve.

```
bitset<N> is; // #include <bitset> -----//87a
int pr[N], primes = 0; -----//87b
void sieve() { -----//87c
  --- is[2] = true; pr[primes++] = 2; -----//87d
  --- for (int i = 3; i < N; i += 2) is[i] = 1; -----//87e
  --- for (int i = 3; i*i < N; i += 2) -----//87f
  --- if (is[i]) -----//880
  --- for (int j = i*i; j < N; j += i) -----//881
  --- is[j]= 0; -----//882
  --- for (int i = 3; i < N; i += 2) -----//883
  --- if (is[i]) -----//884
  --- pr[primes++] = i;} -----//885
```

6.2. Divisor Sieve.

```
int divisors[N]; // initially 0 -----//829
void divisorSieve() { -----//82a
  --- for (int i = 1; i < N; i++) -----//82b
  --- for (int j = i; j < N; j += i) -----//82c
  --- divisors[j]++;} -----//82d
```

6.3. Number/Sum of Divisors. If a number n is prime factorized where $n = p_1^{e_1} \times p_2^{e_2} \times \cdots \times p_k^{e_k}$, where σ_0 is the number of divisors while σ_1 is the sum of divisors:

$$\sum_{d|n} d^k = \sigma_k(n) = \prod \frac{p_i^{k(e_i)+1} - 1}{p_i - 1}$$
$$\text{Product: } \prod_{d|n} d = n^{\frac{\sigma_1(n)}{2}}$$

6.4. Möbius Sieve. The Möbius function μ is the Möbius inverse of e such that $e(n) = \sum_{d|n} \mu(d)$.

```
bitset<N> is; int mu[N]; -----//868
void mobiusSieve() { -----//869
  --- for (int i = 1; i < N; ++i) mu[i] = 1; -----//86a
  --- for (int i = 2; i < N; ++i) if (!is[i]) { -----//86b
  --- for (int j = i; j < N; j += i){ -----//86c
  --- is[j] = 1; -----//86d
  --- mu[j] *= -1; -----//86e
  --- } -----//86f
  --- for (long long j = 1LL*i*i; j < N; j += i*i) -----//870
  --- mu[j] = 0;} -----//871
```

6.5. Möbius Inversion. Given arithmetic functions f and g :

$$g(n) = \sum_{d|n} f(d) \iff f(n) = \sum_{d|n} \mu(d) g\left(\frac{n}{d}\right)$$

6.6. GCD Subset Counting. Count number of subsets $S \subseteq A$ such that $\gcd(S) = g$ (modifiable).

```
int f[MX+1]; // MX is maximum number of array -----//83d
long long gcnt[MX+1]; // gcnt[G]: answer when gcd==G -----//83e
long long C(int f) {return (1LL << f) - 1;} -----//83f
// f: frequency count -----//840
// C(f): # of subsets of f elements (YOU CAN EDIT) -----//841
void gcd_counter(int a[], int n) { -----//842
  --- memset(f, 0, sizeof f); -----//843
  --- memset(gcnt, 0, sizeof gcnt); -----//844
  --- int mx = 0; -----//845
  --- for (int i = 0; i < n; ++i) { -----//846
  --- f[a[i]] += 1; -----//847
  --- mx = max(mx, a[i]); -----//848
  --- } -----//849
  --- for (int i = mx; i >= 1; --i) { -----//84a
  --- int add = f[i]; -----//84b
  --- long long sub = 0; -----//84c
  --- for (int j = 2*i; j <= mx; j += i) { -----//84d
  --- add += f[j]; -----//84e
  --- sub += gcnt[j]; -----//84f
  --- } -----//850
  --- gcnt[i] = C(add) - sub; -----//851
  --- } // Usage: int subsets_with_gcd_1 = gcnt[1]; -----//852
```

6.7. Euler Totient. Counts all integers from 1 to n that are relatively prime to n in $O(\sqrt{n})$ time.

```
LL totient(LL n) { -----//891
  --- if (n <= 1) return 1; -----//892
  --- LL tot = n; -----//893
  --- for (int i = 2; i * i <= n; i++) { -----//894
  --- if (n % i == 0) tot -= tot / i; -----//895
  --- while (n % i == 0) n /= i; -----//896
  --- } -----//897
  --- if (n > 1) tot -= tot / n; -----//898
  --- return tot; } -----//899
```

6.8. Euler Phi Sieve. Sieve version of Euler totient, runs in $O(N \log N)$ time. Note that $n = \sum_{d|n} \varphi(d)$.

```
bitset<N> is; int phi[N]; -----//872
void phiSieve() { -----//873
  --- for (int i = 1; i < N; ++i) phi[i] = i; -----//874
  --- for (int i = 2; i < N; ++i) if (!is[i]) { -----//875
  --- for (int j = i; j < N; j += i) { -----//876
  --- phi[j] -= phi[j] / i; -----//877
  --- is[j] = true; -----//878
  --- }}} -----//879
```

6.9. Extended Euclidean. Assigns x, y such that $ax + by = \gcd(a, b)$ and returns $\gcd(a, b)$.

```
typedef long long LL; -----//82e
typedef pair<LL, LL> PAIR; -----//82f
LL mod(LL x, LL m) { // use this instead of x % m -----//830
  --- if (m == 0) return 0; -----//831
  --- if (m < 0) m *= -1; -----//832
  --- return (x%m + m) % m; // always nonnegative -----//833
} -----//834
```

```
LL extended_euclid(LL a, LL b, LL &x, LL &y) { -----//835
  --- if (b==0) {x = 1; y = 0; return a;} -----//836
  --- LL g = extended_euclid(b, a%b, x, y); -----//837
  --- LL z = x - a/b*y; -----//838
  --- x = y; y = z; return g; -----//839
} -----//83a
```

6.10. Modular Inverse. Find unique x such that $ax \equiv 1 \pmod{m}$. Returns 0 if no unique solution is found. Please use modulo solver for the non-unique case.

```
LL modinv(LL a, LL m) { -----//85e
  --- LL x, y; LL g = extended_euclid(a, m, x, y); -----//85f
  --- if (g == 1 || g == -1) return mod(x * g, m); -----//860
  --- return 0; // 0 if invalid -----//861
} -----//862
```

6.11. Modulo Solver. Solve for values of x for $ax \equiv b \pmod{m}$. Returns $(-1, -1)$ if there is no solution. Returns a pair (x, M) where solution is $x \bmod M$.

```
PAIR modsolver(LL a, LL b, LL m) { -----//863
  --- LL x, y; LL g = extended_euclid(a, m, x, y); -----//864
  --- if (b % g != 0) return PAIR(-1, -1); -----//865
  --- return PAIR(mod(x*b/g, m/g), abs(m/g)); -----//866
} -----//867
```

6.12. Linear Diophantine. Computes integers x and y such that $ax + by = c$, returns $(-1, -1)$ if no solution. Tries to return positive integer answers for x and y if possible.

```
PAIR null(-1, -1); // needs extended euclidean -----//853
PAIR diophantine(LL a, LL b, LL c) { -----//854
  --- if (!a && !b) return c ? null : PAIR(0, 0); -----//855
  --- if (!a) return c % b ? null : PAIR(0, c / b); -----//856
  --- if (!b) return c % a ? null : PAIR(c / a, 0); -----//857
  --- LL x, y; LL g = extended_euclid(a, b, x, y); -----//858
  --- if (c % g) return null; -----//859
  --- y = mod(y * (c/g), a/g); -----//85a
  --- if (y == 0) y += abs(a/g); // prefer positive sol. -----//85b
  --- return PAIR((c - b*y)/a, y); -----//85c
} -----//85d
```

6.13. Chinese Remainder Theorem. Solves linear congruence $x \equiv b_i \pmod{m_i}$. Returns $(-1, -1)$ if there is no solution. Returns a pair (x, M) where solution is $x \bmod M$.

```
PAIR chinese(LL b1, LL m1, LL b2, LL m2) { -----//81b
  --- LL x, y; LL g = extended_euclid(m1, m2, x, y); -----//81c
  --- if (b1 % g != b2 % g) return PAIR(-1, -1); -----//81d
  --- LL M = abs(m1 / g * m2); -----//81e
  --- return PAIR(mod(mod(x*b2*m1+y*b1*m2, M*g)/g,M),M); -----//81f
} -----//820
PAIR chinese_remainder(LL b[], LL m[], int n) { -----//821
  --- PAIR ans(0, 1); -----//822
  --- for (int i = 0; i < n; ++i) { -----//823
  --- ans = chinese(b[i],m[i],ans.first,ans.second); -----//824
  --- if (ans.second == -1) break; -----//825
  --- } -----//826
  --- return ans; -----//827
} -----//828
```


6.13.1. *Super Chinese Remainder.* Solves linear congruence $a_ix \equiv b_i \pmod{m_i}$. Returns $(-1, -1)$ if there is no solution.

```
PAIR super_chinese(LL a[], LL b[], LL m[], int n) {
    PAIR ans(0, 1);
    for (int i = 0; i < n; ++i) {
        PAIR two = modsolver(a[i], b[i], m[i]);
        if (two.second == -1) return two;
        ans = chinese(ans.first, ans.second, two.first, two.second);
        if (ans.second == -1) break;
    }
    return ans;
}
```

7. ALGEBRA

7.1. **Fast Fourier Transform.** Compute the Discrete Fourier Transform (DFT) of a polynomial in $O(n \log n)$ time.

```
struct poly {
    double a, b;
    poly(double a=0, double b=0): a(a), b(b) {}
    poly operator+(const poly& p) const {
        return poly(a + p.a, b + p.b);
    }
    poly operator-(const poly& p) const {
        return poly(a - p.a, b - p.b);
    }
    poly operator*(const poly& p) const {
        return poly(a*p.a - b*p.b, a*p.b + b*p.a);
    }
};

void fft(poly in[], poly p[], int n, int s) {
    if (n < 1) return;
    if (n == 1) {p[0] = in[0]; return;}
    n >>= 1; fft(in, p, n, s << 1);
    fft(in + s, p + n, n, s << 1);
    poly w(1), wn(cos(M_PI/n), sin(M_PI/n));
    for (int i = 0; i < n; ++i) {
        poly even = p[i], odd = p[i + n];
        p[i] = even + w * odd;
        p[i + n] = even - w * odd;
        w = w * wn;
    }
}

void fft(poly p[], int n) {
    poly *f = new poly[n]; fft(p, f, n, 1);
    copy(f, f + n, p); delete[] f;
}

void inverse_fft(poly p[], int n) {
    for(int i=0; i<n; i++) {p[i].b *= -1;} fft(p, n);
    for(int i=0; i<n; i++) {p[i].a/=n; p[i].b/= -1*n;}
}
```

7.2. **FFT Polynomial Multiplication.** Multiply integer polynomials a, b of size an, bn using FFT in $O(n \log n)$. Stores answer in an array c , rounded to the nearest integer (or double).

```
// note: c[] should have size of at least (an+bn)
int mult(int a[],int an,int b[],int bn,int c[]) {
    int n, degree = an + bn - 1;
    for (n = 1; n < degree; n <= 1); // power of 2
```

```
poly *A = new poly[n], *B = new poly[n];
copy(a, a + an, A); fill(A + an, A + n, 0);
copy(b, b + bn, B); fill(B + bn, B + n, 0);
fft(A, n); fft(B, n);
for (int i = 0; i < n; i++) A[i] = A[i] * B[i];
inverse_fft(A, n);
for (int i = 0; i < degree; i++)
    c[i] = int(A[i].a + 0.5); // same as round(A[i].a)
delete[] A, B; return degree;
```

7.3. **Polynomial Long Division.** Divide two polynomials A and B to get Q and R , where $\frac{A}{B} = Q + \frac{R}{B}$.

```
typedef vector<double> Poly;
Poly Q, R; // quotient and remainder
void trim(Poly& A) { // remove trailing zeroes
    while (!A.empty() && abs(A.back()) < EPS)
        A.pop_back();
}

void divide(Poly A, Poly B) {
    if (B.size() == 0) throw exception();
    if (A.size() < B.size()) {Q.clear(); R=A; return;}
    Q.assign(A.size() - B.size() + 1, 0);
    Poly part;
    while (A.size() >= B.size()) {
        int As = A.size(), Bs = B.size();
        part.assign(As, 0);
        for (int i = 0; i < Bs; i++)
            part[As-Bs+i] = B[i];
        double scale = Q[As-Bs] = A[As-1] / part[As-1];
        for (int i = 0; i < As; i++)
            A[i] -= part[i] * scale;
        trim(A);
    }
    R = A; trim(Q);
}
```

7.4. **Matrix Multiplication.** Multiplies matrices $A_{p \times q}$ and $B_{q \times r}$ in $O(n^3)$ time, modulo MOD.

```
long[][] multiply(long A[][], long B[][]) {
    int p = A.length, q = A[0].length, r = B[0].length;
    // if(q != B.length) throw new Exception("((");
    long AB[][] = new long[p][r];
    for (int i = 0; i < p; i++)
        for (int j = 0; j < r; j++)
            for (int k = 0; k < q; k++)
                (AB[i][k] += A[i][j] * B[j][k]) %= MOD;
    return AB;
}
```

7.5. **Matrix Power.** Computes for B^e in $O(n^3 \log e)$ time. Refer to Matrix Multiplication.

```
long[][] power(long B[][], long e) {
    int n = B.length;
    long ans[][] = new long[n][n];
    for (int i = 0; i < n; i++) ans[i][i] = 1;
    while (e > 0) {
        if (e % 2 == 1) ans = multiply(ans, B);
        B = multiply(B, B); e /= 2;
    }
    return ans;
}
```

7.6. **Fibonacci Matrix.** Fast computation for n th Fibonacci $\{F_1, F_2, \dots, F_n\}$ in $O(\log n)$:

$$\begin{bmatrix} F_n \\ F_{n-1} \end{bmatrix} = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}^n \times \begin{bmatrix} F_2 \\ F_1 \end{bmatrix}$$

7.7. **Gauss-Jordan/Matrix Determinant.** Row reduce matrix A in $O(n^3)$ time. Returns true if a solution exists.

```
boolean gaussJordan(double A[][]) {
    int n = A.length, m = A[0].length;
    boolean singular = false;
    // double determinant = 1;
    for (int i=0, p=0; i<n && p<m; i++, p++) {
        for (int k = i + 1; k < n; k++) {
            if (Math.abs(A[k][p]) > EPS) { // swap
                double t=A[i]; A[i]=A[k]; A[k]=t;
                break;
            }
        }
        // determinant *= A[i][p];
        if (Math.abs(A[i][p]) < EPS)
            { singular = true; i--; continue; }
        for (int j = m-1; j >= p; j--) A[i][j]/= A[i][p];
        for (int k = 0; k < n; k++) {
            if (i == k) continue;
            for (int j = m-1; j >= p; j--)
                A[k][j] -= A[k][p] * A[i][j];
        }
    }
    return !singular;
}
```

8. COMBINATORICS

8.1. **Lucas Theorem.** Compute $\binom{n}{k} \pmod p$ in $O(p + \log_p n)$ time, where p is a prime.

```
LL f[P], lid; // P: biggest prime
LL lucas(LL n, LL k, int p) {
    if (k == 0) return 1;
    if (n < p && k < p) {
        if (lid != p) {
            lid = p; f[0] = 1;
            for (int i = 0; i < p; ++i) f[i]=f[i-1]*i%p;
        }
        return f[n] * modpow(f[n-k]*f[k]%p, p-2, p) % p;
    }
    return lucas(n/p,k/p,p) * lucas(n%p,k%p,p) % p;
}
```

8.2. **Granville's Theorem.** Compute $\binom{n}{k} \pmod m$ (for any m) in $O(m^2 \log^2 n)$ time.

```
def fprime(n, p):
    # counts the number of prime divisors of n!
    pk, ans = p, 0
    while pk <= n:
        ans += n // pk
        pk *= p
    return ans

def granville(n, k, p, E):
    # n choose k (mod p^E)
    prime_pow = fprime(n,p)-fprime(k,p)-fprime(n-k,p)
```

```
-- if prime_pow >= E: return 0 -----//092
-- e = E - prime_pow -----//093
-- pe = p ** e -----//094
-- r, f = n - k, [1]*pe -----//095
-- for i in range(1, pe): -----//096
--     x = i -----//097
--     if x % p == 0: -----//098
--         x = 1 -----//099
--     f[i] = f[i-1] * x % pe -----//09a
-- numer, denom, negate, ptr = 1, 1, 0, 0 -----//09b
-- while n: -----//09c
--     if f[-1] != 1 and ptr >= e: -----//09d
--         negate ^= (n&1) ^ (k&1) ^ (r&1) -----//09e
--         numer = numer * f[n%pe] % pe -----//09f
--         denom = denom * f[k%pe] % pe * f[r%pe] % pe -----//0a0
--         n, k, r = n//p, k//p, r//p -----//0a1
--         ptr += 1 -----//0a2
-- ans = numer * modinv(denom, pe) % pe -----//0a3
-- if negate and (p != 2 or e < 3): -----//0a4
--     ans = (pe - ans) % pe -----//0a5
-- return mod(ans * p**prime_pow, p**E) -----//0a6
def choose(n, k, m): # generalized (n choose k) mod m -----//0a7
-- factors, x, p = [], m, 2 -----//0a8
-- while p*p <= x: -----//0a9
--     e = 0 -----//0aa
--     while x % p == 0: -----//0ab
--         e += 1 -----//0ac
--         x //= p -----//0ad
--     if e: factors.append((p, e)) -----//0ae
--     p += 1 -----//0af
-- if x > 1: factors.append((x, 1)) -----//0b0
-- crt_array = [granville(n,k,p,e) for p, e in factors] -----//0b1
-- mod_array = [p**e for p, e in factors] -----//0b2
-- return chinese_remainder(crt_array, mod_array)[0] -----//0b3
```

8.3. **Derangements.** Compute the number of permutations with n elements such that no element is at their original position:

$$D(n) = (n - 1)(D(n - 1) + D(n - 2)) = nD(n - 1) + (-1)^n$$

8.4. **Factoradics.** Convert a permutation of n items to factoradics and vice versa in $O(n \log n)$.

```
// use fenwick tree add, sum, and low code -----//078
typedef long long LL; -----//079
void factoradic(int arr[], int n) { // 0 to n-1 -----//07a
-- for (int i = 0; i <= n; i++) fen[i] = 0; -----//07b
-- for (int i = 1; i < n; i++) add(i, 1); -----//07c
-- for (int i = 0; i < n; i++) { -----//07d
--     int s = sum(arr[i]); -----//07e
--     add(arr[i], -1); arr[i] = s; -----//07f
-- } -----//080
void permute(int arr[], int n) { // factoradic to perm -----//081
-- for (int i = 0; i <= n; i++) fen[i] = 0; -----//082
-- for (int i = 1; i < n; i++) add(i, 1); -----//083
-- for (int i = 0; i < n; i++) { -----//084
--     arr[i] = low(arr[i] - 1); -----//085
--     add(arr[i], -1); -----//086
-- } -----//087
```

8.5. **k th Permutation.** Get the next k th permutation of n items, if exists, using factoradics. All values should be from 0 to $n - 1$. Use factoradics methods as discussed above.

```
bool kth_permutation(int arr[], int n, LL k) { -----//0b4
-- factoradic(arr, n); // values from 0 to n-1 -----//0b5
-- for (int i = n-1; i >= 0 && k > 0; --i){ -----//0b6
--     LL temp = arr[i] + k; -----//0b7
--     arr[i] = temp % (n - i); -----//0b8
--     k = temp / (n - i); -----//0b9
-- } -----//0ba
-- permute(arr, n); -----//0bb
-- return k == 0; } -----//0bc
```

8.6. **Catalan Numbers.**

$$C_n = \frac{1}{n+1} \binom{2n}{n} = \binom{2n}{n} - \binom{2n}{n+1}$$

- (1) The number of non-crossing partitions of an n -element set
- (2) The number of expressions with n pairs of parentheses
- (3) The number of ways $n + 1$ factors can be parenthesized
- (4) The number of full binary trees with $n + 1$ leaves
- (5) The number of monotonic lattice paths of an $n \times n$ grid (5-SAT problem)
- (6) The number of triangulations of a convex polygon with $n + 2$ sides (non-rotational)
- (7) The number of permutations $\{1, \dots, n\}$ without a 3-term increasing subsequence
- (8) The number of ways to form a mountain range with n ups and n downs

8.7. **Stirling Numbers.** s_1 : Count the number of permutations of n elements with k disjoint cycles

s_2 : Count the ways to partition a set of n elements into k nonempty subsets

$$s_1(n, k) = \begin{cases} 1 & n = k = 0 \\ s_1(n - 1, k - 1) - (n - 1)s_1(n - 1, k) & n, k > 0 \\ 0 & \text{elsewhere} \end{cases}$$

$$s_2(n, k) = \begin{cases} 1 & n = k = 0 \\ s_2(n - 1, k - 1) + ks_2(n - 1, k) & n, k > 0 \\ 0 & \text{elsewhere} \end{cases}$$

8.8. **Partition Function.** Pregenerate the number of partitions of positive integer n with n positive addends.

$$p(n, k) = \begin{cases} 1 & n = k = 0 \\ 0 & n < k \\ p(n - 1, k - 1) + p(n - k, k) & n \geq k \end{cases}$$

9. GEOMETRY

```
#include <complex> -----//33c
#define x real() -----//33d
#define y imag() -----//33e
typedef std::complex<double> point; // 2D point only -----//33f
const double PI = acos(-1.0), EPS = 1e-7; -----//340
```

9.1. **Dots and Cross Products.**

```
double dot(point a, point b) -----//38a
- {return a.x * b.x + a.y * b.y;} // + a.z * b.z; -----//38b
double cross(point a, point b) -----//38c
- {return a.x * b.y - a.y * b.x;} -----//38d
double cross(point a, point b, point c) -----//38e
- {return cross(a, b) + cross(b, c) + cross(c, a);} -----//38f
double cross3D(point a, point b) { -----//396
- return point(a.x*b.y - a.y*b.x, a.y*b.z - -----//391
- a.z*b.y, a.z*b.x - a.x*b.z);} -----//392
```

9.2. **Angles and Rotations.**

```
double angle(point a, point b, point c) { -----//2ff
- // angle formed by abc in radians: PI < x <= PI -----//300
- return abs(remainder(arg(a-b) - arg(c-b), 2*PI));} -----//301
point rotate(point p, point a, double d) { -----//302
- //rotate point a about pivot p CCW at d radians -----//303
- return p + (a - p) * point(cos(d), sin(d));} -----//304
```

9.3. **Spherical Coordinates.**

$$\begin{aligned} x &= r \cos \theta \cos \phi & r &= \sqrt{x^2 + y^2 + z^2} \\ y &= r \cos \theta \sin \phi & \theta &= \cos^{-1} x/r \\ z &= r \sin \theta & \phi &= \text{atan2}(y, x) \end{aligned}$$

9.4. **Point Projection.**

```
point proj(point p, point v) { -----//417
- // project point p onto a vector v (2D & 3D) -----//418
- return dot(p, v) / norm(v) * v;} -----//419
point projLine(point p, point a, point b) { -----//41a
- // project point p onto line ab (2D & 3D) -----//41b
- return a + dot(p-a, b-a) / norm(b-a) * (b-a);} -----//41c
point projSeg(point p, point a, point b) { -----//41d
- // project point p onto segment ab (2D & 3D) -----//41e
- double s = dot(p-a, b-a) / norm(b-a); -----//41f
- return a + min(1.0, max(0.0, s)) * (b-a);} -----//420
point projPlane(point p, double a, double b, -----//421
- double c, double d) { -----//422
- // project p onto plane ax+by+cz+d=0 (3D) -----//423
- // same as: o + p - project(p - o, n); -----//424
- double k = -d / (a*a + b*b + c*c); -----//425
- point o(a*k, b*k, c*k), n(a, b, c); -----//426
- point v(p.x-o.x, p.y-o.y, p.z-o.z); -----//427
- double s = dot(v, n) / dot(n, n); -----//428
- return point(o.x + p.x + s * n.x, o.y + -----//429
- p.y + s * n.y, o.z + p.z + s * n.z);} -----//42a
```

9.5. **Great Circle Distance.**

```
double greatCircleDist(double lat1, double long1, -----//393
-- double lat2, double long2, double R) { -----//394
- long1 *= PI / 180; lat1 *= PI / 180; // to radians -----//395
- long2 *= PI / 180; lat2 *= PI / 180; -----//396
- return R*acos(sin(lat1)*sin(lat2) + -----//397
- cos(lat1)*cos(lat2)*cos(abs(long1 - long2)));} -----//398
} -----//399
// another version, using actual (x, y, z) -----//39a
double greatCircleDist(point a, point b) { -----//39b
```

```
- return atan2(abs(cross3D(a, b)), dot3D(a, b)); -----//39c
} -----//39d

9.6. Point/Line/Plane Distances.

double distPtLine(point p, double a, double b, -----//370
- double c) { -----//371
- // dist from point p to line ax+by+c=0 -----//372
- return abs(a*p.x + b*p.y + c) / sqrt(a*a + b*b);} -----//373
double distPtLine(point p, point a, point b) { -----//374
- // dist from point p to line ab -----//375
- return abs((a.y - b.y) * (p.x - a.x) + -----//376
- (b.x - a.x) * (p.y - a.y)) / -----//377
- hypot(a.x - b.x, a.y - b.y);} -----//378
double distPtPlane(point p, double a, double b, -----//379
- double c, double d) { -----//37a
- // distance to 3D plane ax + by + cz + d = 0 -----//37b
- return (a*p.x+b*p.y+c*p.z+d)/sqrt(a*a+b*b+c*c); -----//37c
} /*! // distance between 3D lines AB & CD (untested) -----//37d
double distLine3D(point A,point B,point C,point D){ -----//37e
- point u = B - A, v = D - C, w = A - C; -----//37f
- double a = dot(u, u), b = dot(u, v); -----//380
- double c = dot(v, v), d = dot(u, w); -----//381
- double e = dot(v, w), det = a*c - b*b; -----//382
- double s = det < EPS ? 0.0 : (b*e - c*d) / det; -----//383
- double t = det < EPS -----//384
- ? (b > c ? d/b : e/c) // parallel -----//385
- : (a*e - b*d) / det; -----//386
- point top = A + u * s, bot = w - A - v * t; -----//387
- return dist(top, bot); -----//388
} // dist<EPS: intersection */ -----//389

9.7. Intersections.

9.7.1. Line-Segment Intersection. Get intersection points of 2D
lines/segments  $\overline{ab}$  and  $\overline{cd}$ .

point null(HUGE_VAL, HUGE_VAL); -----//3c9
point line_inter(point a, point b, point c, -----//3ca
- point d, bool seg = false) { -----//3cb
- point ab(b.x - a.x, b.y - a.y); -----//3cc
- point cd(d.x - c.x, d.y - c.y); -----//3cd
- point ac(c.x - a.x, c.y - a.y); -----//3ce
- double D = -cross(ab, cd); // determinant -----//3cf
- double Ds = cross(cd, ac); -----//3d0
- double Dt = cross(ab, ac); -----//3d1
- if (abs(D) < EPS) { // parallel -----//3d2
- if (seg && abs(Ds) < EPS) { // collinear -----//3d3
- point p[] = {a, b, c, d}; -----//3d4
- sort(p, p + 4, [](point a, point b) { -----//3d5
- return a.x < b.x-EPS || -----//3d6
- (dist(a,b) < EPS && a.y < b.y-EPS); -----//3d7
- }); -----//3d8
- return dist(p[1], p[2]) < EPS ? p[1] : null; -----//3d9
- } -----//3da
- return null; -----//3db
- } -----//3dc
- double s = Ds / D, t = Dt / D; -----//3dd
- if (seg && (min(s,t)<-EPS||max(s,t)>1+EPS)) -----//3de
```

```
- return null; -----//3df
- return point(a.x + s * ab.x, a.y + s * ab.y); -----//3e0
}/* double A = cross(d-a, b-a), B = cross(c-a, b-a); -----//3e1
return (B*d - A*c)/(B - A); */ -----//3e2

9.7.2. Circle-Line Intersection. Get intersection points of circle at center
c, radius r, and line  $\overline{ab}$ .

std::vector<point> CL_inter(point c, double r, -----//316
- point a, point b) { -----//317
- point p = projLine(c, a, b); -----//318
- double d = abs(c - p); vector<point> ans; -----//319
- if (d > r + EPS); // none -----//31a
- else if (d > r - EPS) ans.push_back(p); // tangent -----//31b
- else if (d < EPS) { // diameter -----//31c
- point v = r * (b - a) / abs(b - a); -----//31d
- ans.push_back(c + v); -----//31e
- ans.push_back(c - v); -----//31f
- } else { -----//320
- double t = acos(d / r); -----//321
- p = c + (p - c) * r / d; -----//322
- ans.push_back(rotate(c, p, t)); -----//323
- ans.push_back(rotate(c, p, -t)); -----//324
- } return ans; -----//325
} -----//326

9.7.3. Circle-Circle Intersection.

std::vector<point> CC_intersection(point c1, -----//305
- double r1, point c2, double r2) { -----//306
- double d = dist(c1, c2); -----//307
- vector<point> ans; -----//308
- if (d < EPS) { -----//309
- if (abs(r1-r2) < EPS); // inf intersections -----//30a
- } else if (r1 < EPS) { -----//30b
- if (abs(d - r2) < EPS) ans.push_back(c1); -----//30c
- } else { -----//30d
- double s = (r1*r1 + d*d - r2*r2) / (2*r1*d); -----//30e
- double t = acos(max(-1.0, min(1.0, s))); -----//30f
- point mid = c1 + (c2 - c1) * r1 / d; -----//310
- ans.push_back(rotate(c1, mid, t)); -----//311
- if (abs(sin(t)) >= EPS) -----//312
- ans.push_back(rotate(c2, mid, -t)); -----//313
- } return ans; -----//314
} -----//315

9.8. Polygon Areas. Find the area of any 2D polygon given as points
in  $O(n)$ .

double area(point p[], int n) { -----//3fd
- double a = 0; -----//3fe
- for (int i = 0, j = n - 1; i < n; j = i++) -----//3ff
- a += cross(p[i], p[j]); -----//400
- return abs(a) / 2; } -----//401

9.8.1. Triangle Area. Find the area of a triangle using only their lengths.
Lengths must be valid.

double area(double a, double b, double c) { -----//436
- double s = (a + b + c) / 2; -----//437
- return sqrt(s*(s-a)*(s-b)*(s-c)); } -----//438
```

```
Cyclic Quadrilateral Area. Find the area of a cyclic quadrilateral using
only their lengths. A quadrilateral is cyclic if its inner angles sum up to
360°.

double area(double a, double b, double c, double d) { ----//36d
- double s = (a + b + c + d) / 2; -----//36e
- return sqrt((s-a)*(s-b)*(s-c)*(s-d)); } -----//36f

9.9. Polygon Centroid. Get the centroid/center of mass of a polygon
in  $O(m)$ .

point centroid(point p[], int n) { -----//402
- point ans(0, 0); -----//403
- double z = 0; -----//404
- for (int i = 0, j = n - 1; i < n; j = i++) { -----//405
- double cp = cross(p[j], p[i]); -----//406
- ans += (p[j] + p[i]) * cp; -----//407
- z += cp; -----//408
- } return ans / (3 * z); } -----//409

9.10. Convex Hull. Get the convex hull of a set of points using Graham-
Andrew's scan. This sorts the points at  $O(n\log n)$ , then performs the
Monotonic Chain Algorithm at  $O(n)$ .

// counterclockwise hull in p[], returns size of hull ----//341
bool xcmp(const point& a, const point& b) -----//342
- {return a.x < b.x || (a.x == b.x && a.y < b.y);} -----//343
int convex_hull(point p[], int n) { -----//344
- sort(p, p + n, xcmp); if (n <= 1) return n; -----//345
- int k = 0; point *h = new point[2 * n]; -----//346
- double zer = EPS; // -EPS to include collinears -----//347
- for (int i = 0; i < n; h[k++] = p[i++]) -----//348
- while (k >= 2 && cross(h[k-2],h[k-1],p[i]) < zer) ----//349
- --k; -----//34a
- for(int i = n-2, t = k; i >= 0; h[k++] = p[i--]) -----//34b
- while (k > t && cross(h[k-2],h[k-1],p[i]) < zer) ----//34c
- --k; -----//34d
- k -= 1 + (h[0].x==h[1].x&&h[0].y==h[1].y ? 1 : 0); -----//34e
- copy(h, h + k, p); delete[] h; return k; } -----//34f

9.11. Point in Polygon. Check if a point is strictly inside (or on the
border) of a polygon in  $O(n)$ .

bool inPolygon(point q, point p[], int n) { -----//40a
- bool in = false; -----//40b
- for (int i = 0, j = n - 1; i < n; j = i++) -----//40c
- in ^= ((p[i].y > q.y) != (p[j].y > q.y)) && -----//40d
- q.x < (p[j].x - p[i].x) * (q.y - p[i].y) / -----//40e
- (p[j].y - p[i].y) + p[i].x); -----//40f
- return in; } -----//410
bool onPolygon(point q, point p[], int n) { -----//411
- for (int i = 0, j = n - 1; i < n; j = i++) -----//412
- if (abs(dist(p[i], q) + dist(p[j], q) - -----//413
- dist(p[i], p[j])) < EPS) -----//414
- return true; -----//415
- return false; } -----//416
```

9.12. **Cut Polygon by a Line.** Cut polygon by line \overline{ab} to its left in $O(n)$, such that $\angle abp$ is counter-clockwise.

```
vector<point> cut(point p[],int n,point a,point b) { -----//364
- vector<point> poly; -----//365
- for (int i = 0, j = n - 1; i < n; j = i++) { -----//366
--- double c1 = cross(a, b, p[j]); -----//367
--- double c2 = cross(a, b, p[i]); -----//368
--- if (c1 > -EPS) poly.push_back(p[j]); -----//369
--- if (c1 * c2 < -EPS) -----//36a
----- poly.push_back(line_inter(p[j], p[i], a, b)); -----//36b
- } return poly; } -----//36c
```

9.13. **Triangle Centers.**

```
point bary(point A, point B, point C, -----//439
----- double a, double b, double c) { -----//43a
- return (A*a + B*b + C*c) / (a + b + c);} -----//43b
point trilinear(point A, point B, point C, -----//43c
----- double a, double b, double c) { -----//43d
- return bary(A,B,C,abs(B-C)*a, -----//43e
----- abs(C-A)*b,abs(A-B)*c);} -----//43f
point centroid(point A, point B, point C) { -----//440
- return bary(A, B, C, 1, 1, 1);} -----//441
point circumcenter(point A, point B, point C) { -----//442
- double a=norm(B-C), b=norm(C-A), c=norm(A-B); -----//443
- return bary(A,B,C,a*(b+c-a),b*(c+a-b),c*(a+b-c));} -----//444
point orthocenter(point A, point B, point C) { -----//445
- return bary(A,B,C, tan(angle(B,A,C)), -----//446
----- tan(angle(A,B,C)), tan(angle(A,C,B)));} -----//447
point incenter(point A, point B, point C) { -----//448
- return bary(A,B,C,abs(B-C),abs(A-C),abs(A-B));} -----//449
// incircle radius given the side lengths a, b, c -----//44a
double inradius(double a, double b, double c) { -----//44b
- double s = (a + b + c) / 2; -----//44c
- return sqrt(s * (s-a) * (s-b) * (s-c)) / s;} -----//44d
point excenter(point A, point B, point C) { -----//44e
- double a = abs(B-C), b = abs(C-A), c = abs(A-B); -----//44f
- return bary(A, B, C, -a, b, c); -----//450
- // return bary(A, B, C, a, -b, c); -----//451
- // return bary(A, B, C, a, b, -c); -----//452
} -----//453
point brocard(point A, point B, point C) { -----//454
- double a = abs(B-C), b = abs(C-A), c = abs(A-B); -----//455
- return bary(A,B,C,c/b*a,a/c*b,b/a*c); // CCW -----//456
- // return bary(A,B,C,b/c*a,c/a*b,a/b*c); // CW -----//457
} -----//458
point symmedian(point A, point B, point C) { -----//459
- return bary(A,B,C,norm(B-C),norm(C-A),norm(A-B));} -----//45a
```

9.14. **Convex Polygon Intersection.** Get the intersection of two convex polygons in $O(n^2)$.

```
std::vector<point> convex_polygon_inter(point a[], -----//350
--- int an, point b[], int bn) { -----//351
- point ans[an + bn + an*bn]; -----//352
- int size = 0; -----//353
- for (int i = 0; i < an; ++i) -----//354
--- if (inPolygon(a[i],b,bn) || onPolygon(a[i],b,bn)) -----//355
----- ans[size++] = a[i]; -----//356
```

```
- for (int i = 0; i < bn; ++i) -----//357
--- if (inPolygon(b[i],a,an) || onPolygon(b[i],a,an)) -----//358
----- ans[size++] = b[i]; -----//359
- for (int i = 0, I = an - 1; i < an; I = i++) -----//35a
--- for (int j = 0, J = bn - 1; j < bn; J = j++) { -----//35b
----- try { -----//35c
----- point p=line_inter(a[i],a[I],b[j],b[J],true); -----//35d
----- ans[size++] = p; -----//35e
----- } catch (exception ex) {} -----//35f
--- } -----//360
- size = convex_hull(ans, size); -----//361
- return vector<point>(ans, ans + size); -----//362
} -----//363
```

9.15. **Pick's Theorem for Lattice Points.** Count points with integer coordinates inside and on the boundary of a polygon in $O(n)$ using Pick's theorem: $\text{Area} = I + B/2 - 1$.

```
int interior(point p[], int n) -----//3f6
- {return area(p,n) - boundary(p,n) / 2 + 1;} -----//3f7
int boundary(point p[], int n) { -----//3f8
- int ans = 0; -----//3f9
- for (int i = 0, j = n - 1; i < n; j = i++) -----//3fa
--- ans += gcd(p[i].x - p[j].x, p[i].y - p[j].y); -----//3fb
- return ans;} -----//3fc
```

9.16. **Minimum Enclosing Circle.** Get the minimum bounding ball that encloses a set of points (2D or 3D) in Θn .

```
pair<point, double> bounding_ball(point p[], int n){ -----//3e3
- random_shuffle(p, p + n); -----//3e4
- point center(0, 0); double radius = 0; -----//3e5
- for (int i = 0; i < n; ++i) { -----//3e6
--- if (dist(center, p[i]) > radius + EPS) { -----//3e7
----- center = p[i]; radius = 0; -----//3e8
----- for (int j = 0; j < i; ++j) -----//3e9
----- if (dist(center, p[j]) > radius + EPS) { -----//3ea
----- center.x = (p[i].x + p[j].x) / 2; -----//3eb
----- center.y = (p[i].y + p[j].y) / 2; -----//3ec
----- // center.z = (p[i].z + p[j].z) / 2; -----//3ed
----- radius = dist(center, p[i]); // midpoint -----//3ee
----- for (int k = 0; k < j; ++k) -----//3ef
----- if (dist(center, p[k]) > radius + EPS) { -----//3f0
----- center=circumcenter(p[i], p[j], p[k]); -----//3f1
----- radius = dist(center, p[i]); -----//3f2
----- }}}} -----//3f3
- return make_pair(center, radius); -----//3f4
} -----//3f5
```

9.17. **Shamos Algorithm.** Solve for the polygon diameter in $O(n \log n)$.

```
double shamos(point p[], int n) { -----//42b
- point *h = new point[n+1]; copy(p, p + n, h); -----//42c
- int k = convex_hull(h, n); if (k <= 2) return 0; -----//42d
- h[k] = h[0]; double d = HUGE_VAL; -----//42e
- for (int i = 0, j = 1; i < k; ++i) { -----//42f
--- while (distPtLine(h[j+1], h[i], h[i+1]) >= -----//430
----- distPtLine(h[j], h[i], h[i+1])) { -----//431
----- j = (j + 1) % k; -----//432
--- } -----//433
```

```
--- d = min(d, distPtLine(h[j], h[i], h[i+1])); -----//434
- } return d; } -----//435
```

9.18. **kD Tree.** Get the k -nearest neighbors of a point within pruned radius in $O(k \log k \log n)$.

```
#define cpoint const point& -----//39e
bool cmpx(cpoint a, cpoint b) {return a.x < b.x;} -----//39f
bool cmpy(cpoint a, cpoint b) {return a.y < b.y;} -----//3a0
struct KDTree { -----//3a1
- KDTree(point p[],int n): p(p), n(n) {build(0,n);} -----//3a2
- priority_queue< pair<double, point*> > pq; -----//3a3
- point *p; int n, k; double qx, qy, prune; -----//3a4
- void build(int L, int R, bool dvx=false) { -----//3a5
--- if (L >= R) return; -----//3a6
--- int M = (L + R) / 2; -----//3a7
--- nth_element(p + L, p + M, p + R, dvx?cmpx:cmpy); -----//3a8
--- build(L, M, !dvx); build(M + 1, R, !dvx); -----//3a9
- } -----//3aa
- void dfs(int L, int R, bool dvx) { -----//3ab
--- if (L >= R) return; -----//3ac
--- int M = (L + R) / 2; -----//3ad
--- double dx = qx - p[M].x, dy = qy - p[M].y; -----//3ae
--- double delta = dvx ? dx : dy; -----//3af
--- double D = dx * dx + dy * dy; -----//3b0
--- if(D<=prune && (pq.size()<k || D<pq.top().first)){ -----//3b1
----- pq.push(make_pair(D, &p[M])); -----//3b2
--- if (pq.size() > k) pq.pop(); -----//3b3
--- } -----//3b4
--- int nL = L, nR = M, fL = M + 1, fR = R; -----//3b5
--- if (delta > 0) {swap(nL, fL); swap(nR, fR);} -----//3b6
--- dfs(nL, nR, !dvx); -----//3b7
--- D = delta * delta; -----//3b8
--- if (D<=prune && (pq.size()<k || D<pq.top().first)) -----//3b9
--- dfs(fL, fR, !dvx); -----//3ba
- } -----//3bb
- // returns k nearest neighbors of (x, y) in tree -----//3bc
- // usage: vector<point> ans = tree.knn(x, y, 2); -----//3bd
- vector<point> knn(double x, double y, -----//3be
----- int k=1, double r=-1) { -----//3bf
--- qx=x; qy=y; this->k=k; prune=r<0?HUGE_VAL:r*r; -----//3c0
--- dfs(0, n, false); vector<point> v; -----//3c1
--- while (!pq.empty()) { -----//3c2
----- v.push_back(*pq.top().second); -----//3c3
----- pq.pop(); -----//3c4
--- } reverse(v.begin(), v.end()); -----//3c5
--- return v; -----//3c6
- } -----//3c7
}; -----//3c8
```

9.19. **Line Sweep (Closest Pair).** Get the closest pair distance of a set of points in $O(n \log n)$ by sweeping a line and keeping a bounded rectangle. Modifiable for other metrics such as Minkowski and Manhattan distance. For external point queries, see kD Tree.

```
bool cmpy(const point& a, const point& b) -----//327
- {return a.y < b.y;} -----//328
double closest_pair_sweep(point p[], int n) { -----//329
- if (n <= 1) return HUGE_VAL; -----//32a
```



```
- sort(p, p + n, cmpy); -----//32b
- set<point> box; box.insert(p[0]); -----//32c
- double best = 1e13; // infinity, but not HUGE_VAL -----//32d
- for (int L = 0, i = 1; i < n; ++i) { -----//32e
--- while(L < i && p[i].y - p[L].y > best) -----//32f
---- box.erase(p[L++]); -----//330
--- point bound(p[i].x - best, p[i].y - best); -----//331
--- set<point>::iterator it= box.lower_bound(bound); -----//332
--- while (it != box.end() && p[i].x+best >= it->x){ -----//333
---- double dx = p[i].x - it->x; -----//334
---- double dy = p[i].y - it->y; -----//335
---- best = min(best, sqrt(dx*dx + dy*dy)); -----//336
---- ++it; -----//337
--- } -----//338
--- box.insert(p[i]); -----//339
- } return best; -----//33a
} -----//33b
```

9.20. **Line upper/lower envelope.** To find the upper/lower envelope of a collection of lines $a_i + b_ix$, plot the points (b_i, a_i) , add the point $(0, \pm\infty)$ (depending on if upper/lower envelope is desired), and then find the convex hull.

9.21. **Formulas.** Let $a = (a_x, a_y)$ and $b = (b_x, b_y)$ be two-dimensional vectors.

- $a \cdot b = |a||b| \cos \theta$, where θ is the angle between a and b .
- $a \times b = |a||b| \sin \theta$, where θ is the signed angle between a and b .
- $a \times b$ is equal to the area of the parallelogram with two of its sides formed by a and b . Half of that is the area of the triangle formed by a and b .
- The line going through a and b is $Ax + By = C$ where $A = b_y - a_y$, $B = a_x - b_x$, $C = Aa_x + Ba_y$.
- Two lines $A_1x + B_1y = C_1$, $A_2x + B_2y = C_2$ are parallel iff. $D = A_1B_2 - A_2B_1$ is zero. Otherwise their unique intersection is $(B_2C_1 - B_1C_2, A_1C_2 - A_2C_1)/D$.
- **Euler's formula:** $V - E + F = 2$
- Side lengths a, b, c can form a triangle iff. $a + b > c$, $b + c > a$ and $a + c > b$.
- Sum of internal angles of a regular convex n -gon is $(n - 2)\pi$.
- **Law of sines:** $\frac{a}{\sin A} = \frac{b}{\sin B} = \frac{c}{\sin C}$
- **Law of cosines:** $b^2 = a^2 + c^2 - 2ac \cos B$
- Internal tangents of circles $(c_1, r_1), (c_2, r_2)$ intersect at $(c_1r_2 + c_2r_1)/(r_1 + r_2)$, external intersect at $(c_1r_2 - c_2r_1)/(r_1 + r_2)$.

10. OTHER ALGORITHMS

10.1. **Coordinate Compression.**

10.2. **2SAT.**

10.3. **Nth Permutation.**

10.4. **Floyd's Cycle-Finding.**

10.5. **Simulated Annealing.**

10.6. **Hexagonal Grid Algorithms.**

| | | | |
|---|--|---|--|
| 11. USEFUL INFORMATION (CLEAN THIS UP!!) | | · sufficient: QI and $C[b][c] \leq C[a][d], a \leq b \leq c \leq d$ | |
| 12. Misc | | | |
| 12.1. Debugging Tips. | | | |
| <ul style="list-style-type: none">Stack overflow? Recursive DFS on tree that is actually a long path?Floating-point numbers<ul style="list-style-type: none">Getting NaN? Make sure acos etc. are not getting values out of their range (perhaps 1+eps).Rounding negative numbers?Outputting in scientific notation?Wrong Answer?<ul style="list-style-type: none">Read the problem statement again!Are multiple test cases being handled correctly? Try repeating the same test case many times.Integer overflow?Think very carefully about boundaries of all input parametersTry out possible edge cases:<ul style="list-style-type: none">* $n = 0, n = -1, n = 1, n = 2^{31} - 1$ or $n = -2^{31}$* List is empty, or contains a single element* n is even, n is odd* Graph is empty, or contains a single vertex* Graph is a multigraph (loops or multiple edges)* Polygon is concave or non-simpleIs initial condition wrong for small cases?Are you sure the algorithm is correct?Explain your solution to someone.Are you using any functions that you don't completely understand? Maybe STL functions?Maybe you (or someone else) should rewrite the solution?Can the input line be empty?Run-Time Error?<ul style="list-style-type: none">Is it actually Memory Limit Exceeded? | <ul style="list-style-type: none">GreedyRandomizedOptimizations<ul style="list-style-type: none">Use bitset (/64)Switch order of loops (cache locality)Process queries offline<ul style="list-style-type: none">Mo's algorithmSquare-root decompositionPrecomputationEfficient simulation<ul style="list-style-type: none">Mo's algorithmSqrt decompositionStore 2^k jump pointersData structure techniques<ul style="list-style-type: none">Sqrt bucketsStore 2^k jump pointers2^k merging trickCounting<ul style="list-style-type: none">Inclusion-exclusion principleGenerating functionsGraphs<ul style="list-style-type: none">Can we model the problem as a graph?Can we use any properties of the graph?Strongly connected componentsCycles (or odd cycles)Bipartite (no odd cycles)<ul style="list-style-type: none">* Bipartite matching* Hall's marriage theorem* Stable MarriageCut vertex/bridgeBiconnected componentsDegrees of vertices (odd/even)Trees<ul style="list-style-type: none">* Heavy-light decomposition* Centroid decomposition* Least common ancestor* Centers of the treeEulerian path/circuitChinese postman problemTopological sort(Min-Cost) Max FlowMin Cut<ul style="list-style-type: none">* Maximum Density SubgraphHuffman CodingMin-Cost ArborescenceSteiner TreeKirchoff's matrix tree theoremPrüfer sequencesLovász ToggleLook at the DFS tree (which has no cross-edges)Is the graph a DFA or NFA?<ul style="list-style-type: none">* Is it the Synchronizing word problem?Mathematics<ul style="list-style-type: none">Is the function multiplicative?Look for a pattern | <ul style="list-style-type: none">Permutations<ul style="list-style-type: none">* Consider the cycles of the permutationFunctions<ul style="list-style-type: none">* Sum of piecewise-linear functions is a piecewise-linear function* Sum of convex (concave) functions is convex (concave)Modular arithmetic<ul style="list-style-type: none">* Chinese Remainder Theorem* Linear CongruenceSieveSystem of linear equationsValues too big to represent?<ul style="list-style-type: none">* Compute using the logarithm* Divide everything by some large valueLinear programming<ul style="list-style-type: none">* Is the dual problem easier to solve?Can the problem be modeled as a different combinatorial problem? Does that simplify calculations? <ul style="list-style-type: none">Logic<ul style="list-style-type: none">2-SATXOR-SAT (Gauss elimination or Bipartite matching)Meet in the middleOnly work with the smaller half ($\log(n)$)Strings<ul style="list-style-type: none">Trie (maybe over something weird, like bits)Suffix arraySuffix automaton (+DP?)Aho-CorasickeerTreeWork with $S + S$HashingEuler tour, tree to arraySegment trees<ul style="list-style-type: none">Lazy propagationPersistentImplicitSegment tree of XGeometry<ul style="list-style-type: none">Minkowski sum (of convex sets)Rotating calipersSweep line (horizontally or vertically?)Sweep angleConvex hullFix a parameter (possibly the answer).Are there few distinct values?Binary searchSliding Window (+ Monotonic Queue)Computing a Convolution? Fast Fourier TransformComputing a 2D Convolution? FFT on each row, and then on each columnExact Cover (+ Algorithm X)Cycle-FindingWhat is the smallest set of values that identify the solution? The cycle structure of the permutation? The powers of primes in the factorization?Look at the complement problem | |
| 12.2. Solution Ideas. | | | |
| <ul style="list-style-type: none">Dynamic Programming<ul style="list-style-type: none">Parsing CFGs: CYK AlgorithmDrop a parameter, recover from othersSwap answer and a parameterWhen grouping: try splitting in two2^k trickWhen optimizing<ul style="list-style-type: none">* Convex hull optimization<ul style="list-style-type: none">· $dp[i] = \min_{j < i} \{dp[j] + b[j] \times a[i]\}$· $b[j] \geq b[j + 1]$· optionally $a[i] \leq a[i + 1]$· $O(n^2)$ to $O(n)$* Divide and conquer optimization<ul style="list-style-type: none">· $dp[i][j] = \min_{k < j} \{dp[i - 1][k] + C[k][j]\}$· $A[i][j] \leq A[i][j + 1]$· $O(kn^2)$ to $O(kn \log n)$· sufficient: $C[a][c] + C[b][d] \leq C[a][d] + C[b][c], a \leq b \leq c \leq d$ (QI)* Knuth optimization<ul style="list-style-type: none">· $dp[i][j] = \min_{i < k < j} \{dp[i][k] + dp[k][j] + C[i][j]\}$· $A[i][j - 1] \leq A[i][j] \leq A[i + 1][j]$· $O(n^3)$ to $O(n^2)$ | | | |

- Minimize something instead of maximizing
- Immediately enforce necessary conditions. (All values greater than 0? Initialize them all to 1)
- Add large constant to negative numbers to make them positive
- Counting/Bucket sort

13. FORMULAS

- **Legendre symbol:** $(\frac{a}{b}) = a^{(b-1)/2} \pmod{b}$, b odd prime.
- **Heron’s formula:** A triangle with side lengths a, b, c has area $\sqrt{s(s-a)(s-b)(s-c)}$ where $s = \frac{a+b+c}{2}$.
- **Pick’s theorem:** A polygon on an integer grid strictly containing i lattice points and having b lattice points on the boundary has area $i + \frac{b}{2} - 1$. (Nothing similar in higher dimensions)
- **Euler’s totient:** The number of integers less than n that are coprime to n are $n \prod_{p|n} \left(1 - \frac{1}{p}\right)$ where each p is a distinct prime factor of n .
- **König’s theorem:** In any bipartite graph $G = (L \cup R, E)$, the number of edges in a maximum matching is equal to the number of vertices in a minimum vertex cover. Let U be the set of unmatched vertices in L , and Z be the set of vertices that are either in U or are connected to U by an alternating path. Then $K = (L \setminus Z) \cup (R \cap Z)$ is the minimum vertex cover.
- A minumum Steiner tree for n vertices requires at most $n-2$ additional Steiner vertices.
- The number of vertices of a graph is equal to its minimum vertex cover number plus the size of a maximum independent set.
- **Lagrange polynomial** through points $(x_0, y_0), \dots, (x_k, y_k)$ is $L(x) = \sum_{j=0}^k y_j \prod_{\substack{0 \leq m \leq k \\ m \neq j}} \frac{x-x_m}{x_j-x_m}$
- **Hook length formula:** If λ is a Young diagram and $h_\lambda(i, j)$ is the hook-length of cell (i, j) , then then the number of Young tableaux $d_\lambda = n! / \prod h_\lambda(i, j)$.
- **Möbius inversion formula:** If $f(n) = \sum_{d|n} g(d)$, then $g(n) = \sum_{d|n} \mu(d) f(n/d)$. If $f(n) = \sum_{m=1}^n g(\lfloor n/m \rfloor)$, then $g(n) = \sum_{m=1}^n \mu(m) f(\lfloor \frac{n}{m} \rfloor)$.
- #primitive pythagorean triples with hypotenuse $< n$ approx $n/(2\pi)$.
- **Frobenius Number:** largest number which can’t be expressed as a linear combination of numbers a_1, \dots, a_n with non-negative coefficients. $g(a_1, a_2) = a_1 a_2 - a_1 - a_2$, $N(a_1, a_2) = (a_1 - 1)(a_2 - 1)/2$. $g(d \cdot a_1, d \cdot a_2, a_3) = d \cdot g(a_1, a_2, a_3) + a_3(d - 1)$. An integer $x > (\max_i a_i)^2$ can be expressed in such a way iff. $x \mid \gcd(a_1, \dots, a_n)$.

13.1. Physics.

- **Snell’s law:** $\frac{\sin \theta_1}{v_1} = \frac{\sin \theta_2}{v_2}$

13.2. **Markov Chains.** A Markov Chain can be represented as a weighted directed graph of states, where the weight of an edge represents the probability of transitioning over that edge in one timestep. Let $P^{(m)} = (p_{ij}^{(m)})$ be the probability matrix of transitioning from state i to state j in m timesteps, and note that $P^{(1)}$ is the adjacency matrix of the graph. **Chapman-Kolmogorov:** $p_{ij}^{(m+n)} = \sum_k p_{ik}^{(m)} p_{kj}^{(n)}$. It follows that $P^{(m+n)} = P^{(m)} P^{(n)}$ and $P^{(m)} = P^m$. If $p^{(0)}$ is the initial probability distribution (a vector), then $p^{(0)} P^{(m)}$ is the probability distribution after m timesteps.

The return times of a state i is $R_i = \{m \mid p_{ii}^{(m)} > 0\}$, and i is *aperiodic* if $\gcd(R_i) = 1$. A MC is aperiodic if any of its vertices is aperiodic. A MC is *irreducible* if the corresponding graph is strongly connected.

A distribution π is stationary if $\pi P = \pi$. If MC is irreducible then $\pi_i = 1/\mathbb{E}[T_i]$, where T_i is the expected time between two visits at i . π_j/π_i is the expected number of visits at j in between two consecutive visits at i . A MC is *ergodic* if $\lim_{m \rightarrow \infty} p^{(0)} P^m = \pi$. A MC is ergodic iff. it is irreducible and aperiodic.

A MC for a random walk in an undirected weighted graph (un-weighted graph can be made weighted by adding 1-weights) has $p_{uv} = w_{uv} / \sum_x w_{ux}$. If the graph is connected, then $\pi_u = \sum_x w_{ux} / \sum_v \sum_x w_{vx}$. Such a random walk is aperiodic iff. the graph is not bipartite.

An *absorbing* MC is of the form $P = \begin{pmatrix} Q & R \\ 0 & I_r \end{pmatrix}$. Let $N = \sum_{m=0}^\infty Q^m = (I_t - Q)^{-1}$. Then, if starting in state i , the expected number of steps till absorption is the i -th entry in $N1$. If starting in state i , the probability of being absorbed in state j is the (i, j) -th entry of NR . Many problems on MC can be formulated in terms of a system of recurrence relations, and then solved using Gaussian elimination.

13.3. **Burnside’s Lemma.** Let G be a finite group that acts on a set X . For each g in G let X^g denote the set of elements in X that are fixed by g . Then the number of orbits

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$$

$$Z(S_n) = \frac{1}{n} \sum_{l=1}^n a_l Z(S_{n-l})$$

13.4. **Bézout’s identity.** If (x, y) is any solution to $ax + by = d$ (e.g. found by the Extended Euclidean Algorithm), then all solutions are given by

$$\left(x + k \frac{b}{\gcd(a, b)}, y - k \frac{a}{\gcd(a, b)}\right)$$

13.5. Misc.

13.5.1. *Determinants and PM.*

$$\det(A) = \sum_{\sigma \in S_n} \text{sgn}(\sigma) \prod_{i=1}^n a_{i, \sigma(i)}$$

$$\text{perm}(A) = \sum_{\sigma \in S_n} \prod_{i=1}^n a_{i, \sigma(i)}$$

$$\begin{aligned} pf(A) &= \frac{1}{2^n n!} \sum_{\sigma \in S_{2n}} \text{sgn}(\sigma) \prod_{i=1}^n a_{\sigma(2i-1), \sigma(2i)} \\ &= \sum_{M \in \text{PM}(n)} \text{sgn}(M) \prod_{(i,j) \in M} a_{i,j} \end{aligned}$$

13.5.2. *BEST Theorem.* Count directed Eulerian cycles. Number of OST given by Kirchoff’s Theorem (remove r/c with root) $\# \text{OST}(G, r) \cdot \prod_v (d_v - 1)!$

13.5.3. *Primitive Roots.* Only exists when n is $2, 4, p^k, 2p^k$, where p odd prime. Assume n prime. Number of primitive roots $\phi(\phi(n))$ Let g be primitive root. All primitive roots are of the form g^k where $k, \phi(p)$ are coprime.

k -roots: $g^{i \cdot \phi(n)/k}$ for $0 \leq i < k$

13.5.4. *Sum of primes.* For any multiplicative f :

$$S(n, p) = S(n, p-1) - f(p) \cdot (S(n/p, p-1) - S(p-1, p-1))$$

13.5.5. *Floor.*

$$\lfloor \lfloor x/y \rfloor / z \rfloor = \lfloor x/(yz) \rfloor$$

$$x \% y = x - y \lfloor x/y \rfloor$$