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# Nonlinear differential algorithm to compute all the zeros of a generic polynomial

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### I. INTRODUCTION

*Notation.* Hereafter, for definiteness, we always refer to *monic* polynomials of arbitrary order  $N (N \ge 2)$ ,

$$P_N(z; \vec{c}, \underline{x}) = z^N + \sum_{m=1}^N (c_m z^{N-m}) = \prod_{n=1}^N (z - x_n) , \qquad (1)$$

the *complex* variable z is the argument of the polynomial, indices such as  $n, m, \ell$  run from 1 to N (unless otherwise indicated, see below), the N-vector  $\vec{c}$  has the N coefficients  $c_m$  of the polynomial (1) as its N components,  $\underline{x}$  is the *unordered* set of the N zeros  $x_n$  of the polynomial (1), and we assume all variables to be *complex* (unless otherwise explicitly indicated, see below). We call *generic* any polynomial the *coefficients* and zeros of which are generic complex numbers, and in particular feature zeros which are all different among themselves,  $x_n \neq x_m$  if  $n \neq m$ . Note that the notation  $P_N\left(z;\vec{c},\underline{x}\right)$  is somewhat redundant, since this monic polynomial can be identified by assigning either its N coefficients or its N zeros; indeed the N coefficients  $c_m$  can be expressed in terms of the N zeros  $x_n$  via the following standard formula:

$$c_m = (-1)^m \sum_{\substack{n_1 > n_2 > \dots > n_m = 1}}^N (x_{n_1} x_{n_2} \cdots x_{n_m}), \qquad (2a)$$

so that

$$c_{1} = -(x_{1} + x_{2} + \dots + x_{N}),$$

$$c_{2} = (x_{1} x_{2} + x_{1} x_{3} + \dots + x_{1} x_{N})$$

$$+(x_{2} x_{3} + x_{2} x_{4} + \dots + x_{2} x_{N}) + \dots$$

$$+(x_{N-2} x_{N-1} + x_{N-2} x_{N}) + x_{N-1} x_{N}$$
(2c)

and so on. On the other hand, while the assignment of the *N* coefficients  $c_m$  determines the *N* zeros  $x_n$ —uniquely, up to permutations—of course explicit formulas to accomplish generally this task only exist for  $N \le 4$ .

The investigation of the properties—and of techniques for the numerical computation—of the N zeros  $x_n$  of a polynomial of degree N defined via the assignment of its N coefficients  $c_m$  (see (1)) is a problem that has engaged mathematicians since time immemorial. In this paper, a simple nonlinear differential algorithm suitable to compute numerically *all* the N zeros of a *generic* polynomial of arbitrary degree N is described; I was unable to find a previous description of this algorithm in the literature, but I am aware that my search has not been—indeed, it could not have been—quite complete. This algorithm is described in Section II and proven in Section III.

#### II. RESULTS

It is now convenient to introduce an additional independent variable t, which is hereafter assumed to be *real* and might be interpreted as *time*. Hence the above notation is now extended by writing, in addition to (1), the analogous formula,

$$p_N\left(z;\vec{\gamma}(t),\underline{y}(t)\right) = z^N + \sum_{m=1}^N \left[\gamma_m(t) \ z^{N-m}\right] = \prod_{n=1}^N \left[z - y_n(t)\right],\tag{3}$$

to which notational comments quite analogous to those reported above apply.

There holds then the following.

*Proposition.* Consider the following system of N nonlinear first-order differential equations satisfied by the N zeros  $y_n(t)$  of the polynomial (3),

$$\dot{y}_n(t) = -g(t) \left\{ \prod_{\ell=1, \ \ell \neq n}^{N} \left[ y_n(t) - y_\ell(t) \right]^{-1} \right\} \sum_{m=1}^{N} \left\{ \left[ c_m - \gamma_m(0) \right] \left[ y_n(t) \right]^{N-m} \right\} , \tag{4a}$$

$$g(t) = \frac{\dot{f}(t)}{f(T) - f(0)}$$
 implying  $\int_0^T dt \ g(t) = 1$ . (4b)

Here and below a superimposed dot denote a *t*-differentiation, while the *coefficients*  $c_m$  are those of the polynomial  $P_N\left(z;\vec{c},\underline{x}\right)$ , see (1), the *zeros*  $x_n$  of which we seek, and  $\gamma_m\left(0\right)$  are the *N* coefficients of the polynomial  $p_N\left(z;\vec{\gamma}\left(t\right),y\left(t\right)\right)$ , see (3), at t=0, hence they are related to the "initial" values  $y_n\left(0\right)$  of the *zeros* of this polynomial by the formula (analogous to (2)),

$$\gamma_m(0) = (-1)^m \sum_{n_1 > n_2 > \dots > n_m = 1}^N \left[ y_{n_1}(0) \ y_{n_2}(0) \cdots y_{n_m}(0) \right]. \tag{4c}$$

As for the function f(t), and the *positive* parameter T, see (4b), they can both be assigned essentially *arbitrarily*; but of course so that the function g(t) be finite for  $0 \le t \le T$  hence feature the property displayed by the second equality (4b).

Then

$$x_n = y_n(T) . (4d)$$

It is thus seen that the zeros  $x_n$  of the polynomial  $P_N(z; \vec{c}, \underline{x})$ , see (1), can be computed—once the N coefficients  $c_m$  of this polynomial have been assigned—via the following procedure. Step one: choose (arbitrarily!) N complex numbers  $y_n(0)$ . Step two: compute, via the formulas (4c), the N quantities  $\gamma_m(0)$ . Step three: integrate (numerically) the system of differential equations (4a) from t=0 to t=T, starting from the N initial data  $y_n(0)$ , getting thereby the N values  $y_n(T)$ , which give the sought result, see (4d).

Will this procedure always work? The only possible snag is that the solution  $\vec{y}(t)$  of the "dynamical system" (4a) runs into a singularity during its evolution from t=0 to t=T. The only mechanism whereby this might occur is because during this evolution two different coordinates  $y_n(t)$  might coincide,  $y_\ell(t) = y_n(t)$  for  $\ell \neq n$ , at some value of the *real* variable t in the interval 0 < t < 1 causing the right-hand side of (4a) to blow up. This "collision" might indeed happen, but it is *not* a *generic* phenomenon; hence it will be enough to change the assignment of the (arbitrary!) initial data  $y_n(0)$  to avoid this difficulty; note however that this suggests that to apply this method it will be advisable to always start with *complex* initial data  $y_n(0)$ , even in the case of *real* polynomials with *real* zeros. And note moreover that the numerical integration of the differential equations (4a) with different initial data  $\vec{y}(0)$  and different assignments of the function f(t) and of the parameter T—for instance

$$f(t) = t$$
,  $T = 1$ , implying  $g(t) = 1$ , (5a)

or

$$f(t) = a t^{\lambda}, \lambda > 0, T = 1, \text{ implying } g(t) = \lambda t^{\lambda - 1},$$
 (5b)

or

$$f(t) = \frac{\exp(\lambda t) - 1}{\exp(\lambda T) - 1}, \ \lambda > 0, \ \text{ implying } g(t) = \frac{\lambda \exp(\lambda t)}{\exp(\lambda T) - 1}$$
 (5c)

—allows to assess the *accuracy* of the computation, by comparing the results obtained starting from different assignments of these input data.

Remark. It is plain that this procedure will work more efficiently the closer the, arbitrarily chosen, initial values  $y_n(0)$  are to the N zeros  $x_n$  the values of which one is trying to compute; indeed if the N initial values  $y_n(0)$  happened to coincide with the N zeros  $x_n$ ,  $y_n(0) = x_n$ , this would imply  $y_m(0) = c_m$  (compare (2) with (4c)) hence the right-hand side of the differential equations (4a) would vanish identically, entailing  $\dot{y}_n = 0$  hence  $y_n(1) = y_n(0) = x_n$ , consistently with (4d). This suggests that it might be convenient to employ this technique iteratively, using the output y(T) of each application of it as input y(0) for a subsequent application.

Let us also emphasize that the dependence (via (4c)) of the right-hand sides of the differential equations (4a) upon the initial values  $y_n(0)$  of the dependent variables  $y_n(t)$  implies that these differential equations are rather differential functional equations than ordinary differential equations; but this fact has hardly any relevance on *step three* of the procedure, see above.

Preliminary computations performed by younger colleagues confirm these findings, but a comparison of the actual effectiveness of this technique with that of other methods to compute *all* the *N* zeros of a *generic* polynomial of arbitrary degree *N* is beyond the scope of this short communication, and in any case it is a task to be rather pursued by specialists in numerical analysis if they consider it worthy of their attention.

#### III. PROOF

The proof of the above *Proposition* is actually quite easy (raising thereby some doubts on the novelty of this finding). The starting point is the *identity*,

$$\dot{y}_n(t) = -\left\{ \prod_{\ell=1, \ \ell \neq n}^{N} \left[ y_n(t) - y_\ell(t) \right]^{-1} \right\} \sum_{m=1}^{N} \left\{ \dot{\gamma}_m(t) \left[ y_n(t) \right]^{N-m} \right\} , \tag{6}$$

valid for any t-dependent polynomial with zeros  $y_n(t)$  and coefficients  $\gamma_m(t)$ , see (3); for a proof of this formula see (if need be) Ref. 1. Now make the assignment

$$\gamma_m(t) = \gamma_m(0) + \left[ \frac{f(t) - f(0)}{f(T) - f(0)} \right] [c_m - \gamma_m(0)]$$
 (7a)

consistent with the initial assignment at t = 0 and clearly implying (see (4b))

$$\dot{\gamma}_m(t) = g(t) \left[ c_m - \gamma_m(0) \right] \tag{7b}$$

and

$$\gamma_m(T) = c_m \,. \tag{7c}$$

The insertion of the first of these two formulas, (7b), in (6) yields (4a); while the second, (7c), implies that, at t = T, the polynomial  $p_N(z; \vec{\gamma}(t), \underline{y}(t))$ , see (3), coincides with the polynomial  $P_N(z; \vec{c}, \underline{x})$ , see (1), hence the validity of (4d). Q. E. D.

<sup>&</sup>lt;sup>1</sup> F. Calogero, "New solvable variants of the goldfish many-body problem," Stud. Appl. Math. 137, 123 (2015).