

Knowledge Representation and Reasoning

Exercise Session 5

→ T 0100
→ T 0110
→ T 1110
→ T 1111

T 0001
T 1001
T 0011
T 1111

FINAL

Exercise 1. Type Graph

Let $\varphi = x \mathcal{U} \neg y$

1. Find the types of φ
2. Construct the type graph
3. Identify initial and final types

$0(x \mathcal{U} 1y)$

0/1
0
1

$x \mathcal{U} 1y$

0
0
1

x

0
1
1

y

0
1
1

(*)

Exercise 2. Model Counting

(*)

How many temporal models of length 2 satisfy the formula φ from Exercise 1?

Exercise 3. KR 1

(*)

1. Construct an LTL_f formula describing the following specification of a (simplified) traffic light; you can use the abbreviations seen during the lecture.
 - the light is either green or red, but never both
 - whenever the light is red, it will eventually turn green

Hint: use the propositional variables **green** and **red**.

2. What characterises the **last** timepoint of all models satisfying this specification?

Exercise 4. KR 2

(**)

1. Extend the specification from Exercise 3 to include two traffic lights (with variables $green_i$ and red_i ($i = 1, 2$)) such that the two green lights are never simultaneously on.
2. Is this specification satisfiable? If yes, give a temporal model satisfying it; if not, envision a way to fix it

Exercise 5. Model Size 1

(**)

Build a formula that is satisfied by models of **even** length only, or argue why it cannot exist.

Exercise 6. Model Size 2

(***)

Build a formula that is satisfied by models of **prime** length only, or argue why it cannot exist.

WHEN THE
INIT
FORMULA
IS 1

- 8 NODES



PUT THE 
AROUND

ADD NARROW ↓
TO INITIAL

Exercise 2. Model Counting

(*)

How many temporal models of length 2 satisfy the formula φ from Exercise 1?

$x \cup y$

LENGTH OF A TEMP. MODEL:

HAS $N + 1$ VALUATION
THE LENGTH IS N

0 or 1

SEQUENCE OF
PROPOSITIONAL
VALUATION

$v_0, v_1, v_2, \dots, v_N$

TOT. VALUATION
IS $v_N + 1$ (BECAUSE OF
 v_0)

$x \cup y$

LENGTH $(v_0, v_1, v_2, \dots, v_N) = N$

$v \{x, y\} = \{0, 1\}$

2 cases

| | v_0 | v_1 | v_2 |
|-----------------|---------------|---------------------------------------|---------------------------------------|
| x $\neg x$ | $\neg y$ • | $x, \neg x$ $y, \neg y$ 4 cases | $x, \neg x$ $y, \neg y$ 4 cases |
| x, y | | $\neg y$ $x, \neg y$ 2 cases | 4 cases |
| x, y | | $\neg y$ $x, \neg y$ 2 cases | $\neg y$ $x, \neg x$ 2 cases |

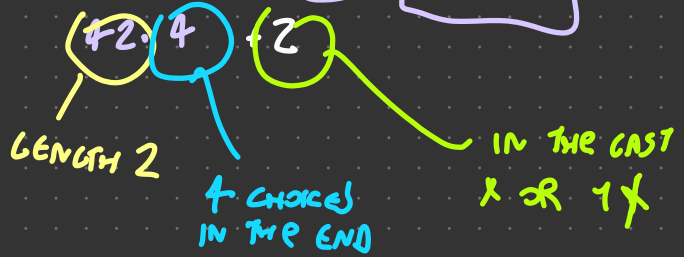
DON'T CARE WHAT HAPPENS NEXT

SAT REGARDLESS OF x TRUE OR NOT

TOT 32 CASES
WHERE $v_0 \models \neg y$
8 CASES OF TEMPORAL MODEL

$$32 + 8 + 2 = \boxed{42} \text{ TEMP. MODEL}$$

$$\text{of LENGTH } \boxed{3} = \boxed{170}$$



Exercise 3. KR 1

(*)

1. Construct an LTL_f formula describing the following specification of a (simplified) traffic light; you can use the abbreviations seen during the lecture.

- 1) • the light is either green or red, but never both
- 2) • whenever the light is red, it will eventually turn green

Hint: use the propositional variables **green** and **red**.

2. What characterises the **last** timepoint of all models satisfying this specification?

GREEN XOR RED

GREEN \leftrightarrow \neg RED

\square = ALWAYS

\diamond = EVENTUALLY

$(GREEN \vee RED) \wedge \neg (GREEN \wedge RED)$

1) $\square (GREEN \text{ XOR } RED)$

2) $\square (RED \rightarrow (RED \cup GREEN))$

$\square (RED \rightarrow \diamond GREEN)$

3) $\square (GREEN_1 \text{ XOR } GREEN_2)$

$\neg (\neg (GREEN_1 \wedge GREEN_2))$

Exercise 4. KR 2

(**)

1. Extend the specification from Exercise 3 to include two traffic lights (with variables $green_i$ and red_i ($i = 1, 2$) such that the two green lights are never simultaneously on.
2. Is this specification satisfiable? If yes, give a temporal model satisfying it; if not, envision a way to fix it

UNSAT

TIME (1) IS GREEN
THE (2) IS RED

NEED TO HAVE
ANOTHER POINT
IN TIME WHEN
IS GREEN

AND SO ON
W/ (1)

IT HAS NO
FINITE MODEL

How to fix it

- NONE IS ON
- BOTH YELLOW??
- BOUNDED TIME
- $\Box (red \rightarrow \Diamond (green \vee ol))$ ∞ -TIMES

NO NEXT POINT
IN TIME

Exercise 5. Model Size 1

(**)

Build a formula that is satisfied by models of **even** length only, or argue why it cannot exist.



$\bar{x} \quad \neg x \quad x \quad \neg x \quad x$

CAN'T END WITH $\neg x$

$$x \wedge \boxed{\neg (x \rightarrow 0x)} \wedge \boxed{(x \rightarrow (\neg 1x \vee 0x))}$$

↑ FIRST ELEMENT IS x

↑ FOR ALL

||

$$\boxed{((x \wedge 0T) \rightarrow 01x)}$$

FOR ODD PUT

$\neg x$ INSTEAD OF x

Exercise 6. Model Size 2

(***)

Build a formula that is satisfied by models of **prime** length only, or argue why it cannot exist.

1, 2, 3, 5, 7, 11, 13

NO EXIST

NOT UNIFORMLY DISTRIBUTED

ALSO THERE ARE INFINITELY MANY PRIME NUMBERS } BUT THE FORMULA IS FINITE

BY CONTRADICTION:

G_φ IS FINITE \rightarrow HAS N TYPES

$p > N$

\uparrow
PRIME NUMBER

THERE IS A PATH OF LENGTH p

$t_0, t_1, t_2, \dots, t_i$ $i < 5$

$t_i = t_5$

