Knowledge Representation and Reasoning

Exercise Session 3

Exercise 1. Subsumption

(*)

Use the **homomorphism method** to verify whether the following subsumption relations hold:

- 1. $A \sqcap C \sqcap \exists r.B \sqsubseteq C \sqcap \exists r.\top$
- 2. $\exists r.(B \sqcap \exists s.\top) \sqsubseteq \exists s.\exists s.\top$
- 3. $\exists r.(B \sqcap \exists s.\top) \sqsubseteq \exists r.\exists s.B$
- 4. $\exists r.(B \sqcap \exists s.A \sqcap \exists s.A)) \sqsubseteq \exists s.\exists s.A$
- 5. $B \cap \exists r.A \cap \exists r.B \cap \exists r.C \sqsubseteq \exists r.(A \cap C) \cap \exists r.(B \cap C)$

Exercise 2. Couter-Models

(*)

For the following pairs of concepts C, D, find an interpretation \mathcal{I} such that $C^{\mathcal{I}} \not\subseteq D^{\mathcal{I}}$

- 1. $C = \exists r. \top, D = \exists r. A$
- 2. $C = A \sqcap C \sqcap \exists r.(A \sqcap B), D = \exists r.(A \sqcap C)$
- 3. $C = \exists r.A \cap \exists r.B, D = \exists r.(A \cap B)$
- 4. $C = \exists r.(A \sqcap B), D = \exists r.A \sqcap \exists s.B$
- 5. $C = A \sqcap C \sqcap \exists r.(A \sqcap B), D = \exists r.(A \sqcap \exists s.\bot)$

Exercise 3. TBox Models

(*)

Construct a model of the following TBox which has at least three elements

$$\{\exists r. (A \sqcap B) \sqsubseteq A \sqcap \exists r. C, \quad C \sqcap \exists s. \top \sqsubseteq B \sqcap \exists s. B, \quad B \sqcap D \sqsubseteq \bot, \quad \exists s. D \sqsubseteq A \sqcap C\}$$

Exercise 4. Disjoint Unions

(***)

Recall the notion of **disjoint union** of interpretations from the previous exercise session. Is the disjoint union of two models of an \mathcal{EL}_{\perp} TBox \mathcal{T} also a model of \mathcal{T} ? **Justify**

Exercise 5. Model Size

(***)

Construct an \mathcal{EL}_{\perp} TBox that has only models with **at most** 5 elements, or argue why it cannot exist.

Exercise 6. Normalisation

(*)

Transform the TBox from Exercise 3 to normal form.

Exercise 7. Reasoning

(*)

Let \mathcal{T} be the TBox from Exercise 3.

- 1. Apply the completion algorithm to check whether the following consequences hold:
 - $\bullet \ \exists r. \exists s. D \sqsubseteq A \sqcap \exists r. \exists s. B$
 - $\bullet \ \ D \sqcap \exists s.D \sqsubseteq B \sqcap \exists r. \top$
 - $\bullet \ \ B \sqcap \exists r. \top \sqsubseteq D \sqcap \exists s. D$
- 2. Construct eventual countermodels

Exercise 8. Completeness

(***)

In the lecture we showed how to build an interpretation \mathcal{I} from the data structure resulting from the completion algorithm, and **claimed it was a model**. As a step towards proving this claim, show that if the TBox contains the GCI $X \sqsubseteq Y$, then the interpretation \mathcal{I} satisfies this GCI.

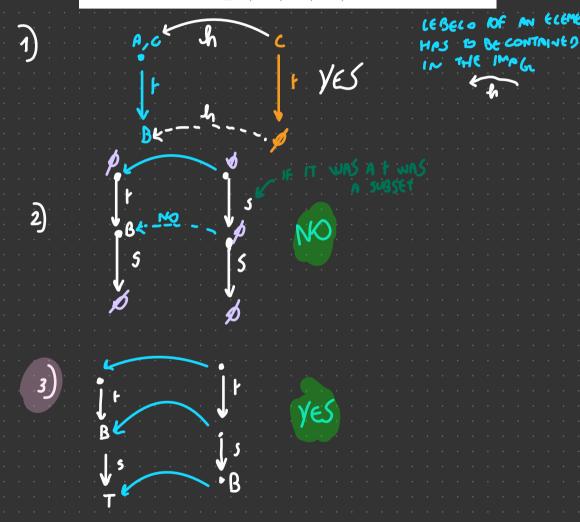
Exercise 9. Inverse Roles

(***)

Using **inverse roles** build a TBox that expresses the knowledge that *humans can only have human children*.

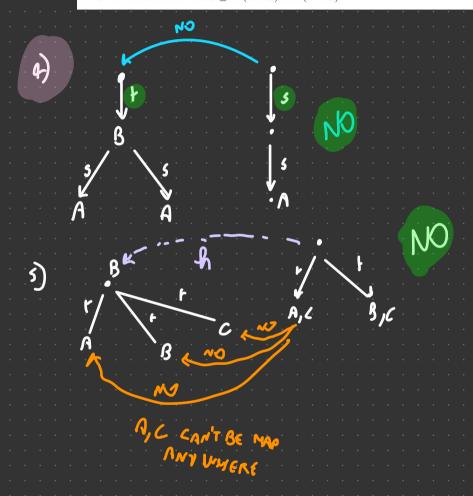
Use the **homomorphism method** to verify whether the following subsumption relations hold:

- 1. $A \sqcap C \sqcap \exists r.B \sqsubseteq C \sqcap \exists r. \top$
- 2. $\exists r.(B \sqcap \exists s.\top) \sqsubseteq \exists s.\exists s.\top$
- 3. $\exists r.(B \sqcap \exists s.\top) \sqsubseteq \exists r.\exists s.B$
- 4. $\exists r.(B \sqcap \exists s.A \sqcap \exists s.A)) \sqsubseteq \exists s.\exists s.A$
- 5. $B \sqcap \exists r.A \sqcap \exists r.B \sqcap \exists r.C \sqsubseteq \exists r.(A \sqcap C) \sqcap \exists r.(B \sqcap C)$



Use the ${f homomorphism}$ ${f method}$ to verify whether the following subsumption relations hold:

- 1. $A \sqcap C \sqcap \exists r.B \sqsubseteq C \sqcap \exists r.\top$
- 2. $\exists r.(B \sqcap \exists s.\top) \sqsubseteq \exists s.\exists s.\top$
- 3. $\exists r.(B \sqcap \exists s.\top) \sqsubseteq \exists r.\exists s.B$
- $4. \ \exists r.(B \sqcap \exists s.A \sqcap \exists s.A) \ \sqsubseteq \exists s.\exists s.A$
- 5. $B \sqcap \exists r.A \sqcap \exists r.B \sqcap \exists r.C \sqsubseteq \exists r.(A \sqcap C) \sqcap \exists r.(B \sqcap C)$



Exercise 2. Couter-Models

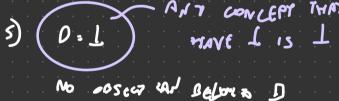
For the following pairs of concepts C, D, find an interpretation \mathcal{I} such that $C^{\mathcal{I}} \not\subseteq D^{\mathcal{I}}$

- 1. $C = \exists r. \top, D = \exists r. A$
- 2. $C = A \sqcap C \sqcap \exists r.(A \sqcap B), D = \exists r.(A \sqcap C)$
- 3. $C = \exists r.A \cap \exists r.B, D = \exists r.(A \cap B)$
- 4. $C = \exists r.(A \sqcap B), D = \exists r.A \sqcap \exists s.B$
- 5. $C = A \sqcap C \sqcap \exists r.(A \sqcap B), D = \exists r.(A \sqcap \exists s.\bot)$

I BE CI WI SECT AND

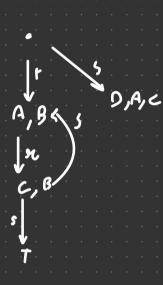
FIND An

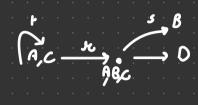
BEONG



Construct a model of the following TBox which has at least three elements

 $\{\exists r. (A \sqcap B) \sqsubseteq A \sqcap \exists r. C, \quad C \sqcap \exists s. \top \sqsubseteq B \sqcap \exists s. B, \quad B \sqcap D \sqsubseteq \bot, \quad \exists s. D \sqsubseteq A \sqcap C\}$



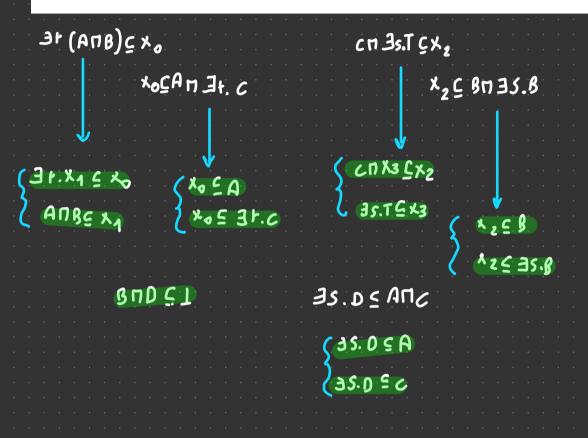


Exercise 3. TBox Models

(*)

Construct a model of the following TBox which has at least three elements

 $\{\exists r.(A \sqcap B) \sqsubseteq A \sqcap \exists r.C, \quad C \sqcap \exists s.\top \sqsubseteq B \sqcap \exists s.B, \quad B \sqcap D \sqsubseteq \bot, \quad \exists s.D \sqsubseteq A \sqcap C\}$



Recall the notion of **disjoint union** of interpretations from the previous exercise session.

Is the disjoint union of two models of an \mathcal{EL}_{\perp} TBox \mathcal{T} also a model of \mathcal{T} ? Justify

$$I = (\Delta^{1}, \cdot^{1}) \quad \Delta^{1} \cap \Delta^{2} = \emptyset$$

$$I = (\Delta^{1}, \cdot^{2}) \quad I + I = (\Delta^{1} \cup \Delta^{2}, \cdot^{1})$$

Exercise 5. Model Size

(***)

Construct an \mathcal{EL}_{\perp} TBox that has only models with **at most** 5 elements, or argue why it cannot exist.

NO 'I CANNOT EXIST

LAN IAVE INFINITE ELEMENT

IF YOU ADD ANOTHER ELEMENT THAT IS OUTSIDE

THE CONSTRAIN OF THE MODEL THE MODEL FTICE

EXISTS AND HAVE ONE ELEMENT MORE

CAN'T PORCE TWO ELEMENTS TO BE COMMECTED

Exercise 6. Normalisation

(*)

Transform the TBox from Exercise 3 to normal form.

Let \mathcal{T} be the TBox from Exercise 3.

- 1. Apply the completion algorithm to check whether the following consequences hold:
 - $\bullet \ \exists r. \exists s. D \sqsubseteq A \sqcap \exists r. \exists s. B$
 - $D \sqcap \exists s.D \sqsubseteq B \sqcap \exists r.\top$
 - $\bullet \ \ B \sqcap \exists r. \top \sqsubseteq D \sqcap \exists s. D$
- 2. Construct eventual countermodels

31.X1 5 % ANBE X1	ASA BSB CSC OSD X15×1 \$25×2		ACT BCT CCT DCT 1151 1251	
800 <u>51</u>				
2x2 35.TEX3	X 25		* 2 ?	
35.D = C 12 = B 12 = 35.B				

Exercise 8. Completeness

(***)

In the lecture we showed how to build an interpretation \mathcal{I} from the data structure resulting from the completion algorithm, and **claimed it was a model**. As a step towards proving this claim, show that if the TBox contains the GCI $X \sqsubseteq Y$, then the interpretation \mathcal{I} satisfies this GCI.

