

Video Lecture 1 Non-ideal Behaviour

ENGN 4213/6213

Digital Systems and Microprocessors



What's in this lecture

- The operation of CMOS logic in nonideal conditions or under a load
- Timing and timing hazards



Resources

- Wakerly (5th edition) 14.3 and 14.4 for static and dynamic electrical behaviour of CMOS (we do not get into as much detail as the book does)
- Wakerly 4.2 on circuit timing and 3.4 on timing hazards

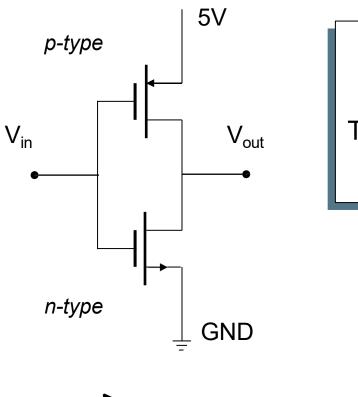


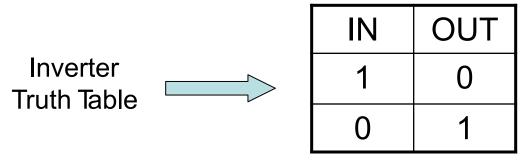
On MOSFET/CMOS again

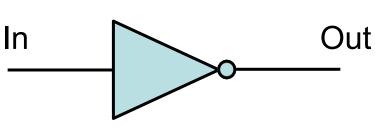
- We remember the following key characteristics
 - MOSFET transistors, when conducting, either behave as a resistor ("ohmic" or "triode" mode) or a current source (saturation mode), depending on the value of V_{GS} w.r.t. V_{DS}
 - CMOS technology combines n-MOS and p-MOS transistors
 - CMOS components use very little power if no logic switching occurs
 - CMOS logic relies on input/output signals being very close to the upper/lower reference voltages (e.g., 0V/5V)

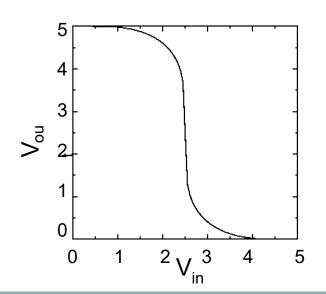


The CMOS inverter





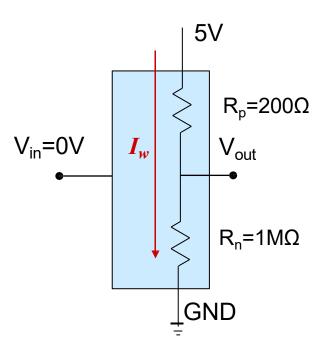






CMOS with nonideal input

Normal operation

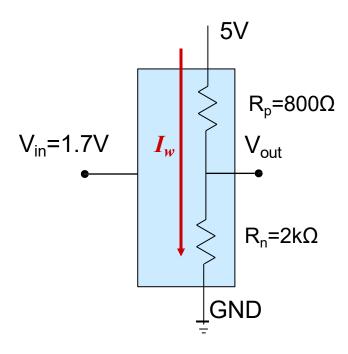


$$I_{w} = 5V/(R_{p}+R_{n}) = 5V/(1.0002M\Omega) < 1\mu A$$

$$P_{w} = 5V \cdot I_{w} = 5\mu W$$

$$V_{out} = 5V - (200\Omega \cdot 1\mu A) = 4.9998V$$

Nonideal input



$$I_w = 5\text{V}/(2.8\text{k}\Omega) = 1.79\text{mA}$$

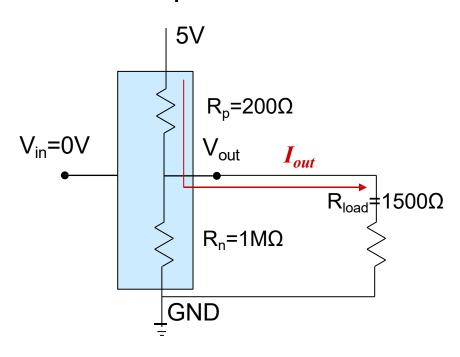
$$P_w = 5\text{V} \cdot I_w = 8.93\text{mW}$$

$$V_{out} = 5V - (800\Omega \cdot 1.79 \text{mA}) = 3.57V$$



CMOS with resistive load

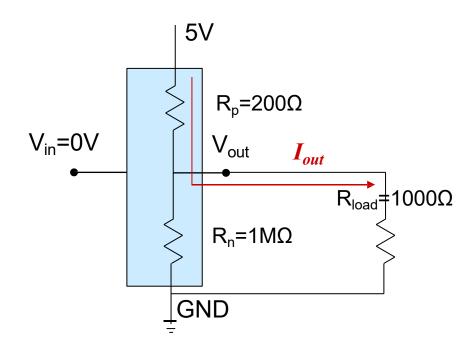
Acceptable load



$$I_{out} \approx 5\text{V}/(\text{R}_{\text{p}} + \text{R}_{\text{load}}) = 5\text{V}/(1.7\text{k}\Omega) = 3\text{mA}$$

 $V_{out} = 5\text{V} - (200\Omega \cdot 3\text{mA}) = 4.4\text{V}$

Unacceptable load



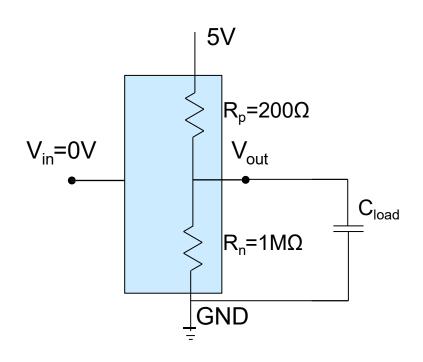
$$I_{out} \approx 5 \text{V}/(\text{R}_{\text{p}} + \text{R}_{\text{load}}) = 5 \text{V}/(1.2 \text{k}\Omega) = 4.16 \text{mA}$$

 $V_{out} = 5 \text{V} - (200\Omega \cdot 4.16 \text{mA}) \neq 4.17 \text{V}$



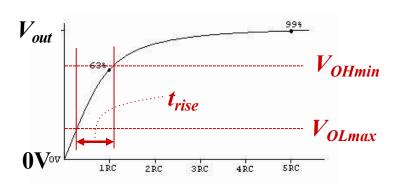
CMOS with a capacitive load

- Capacitive loads are always present.
 - Even the input terminal of a downstream logic stage has capacitance



The charging/discharging of the load capacitance is not instantaneous.

The behaviour can be approximated as a simple RC network with a rise/fall time



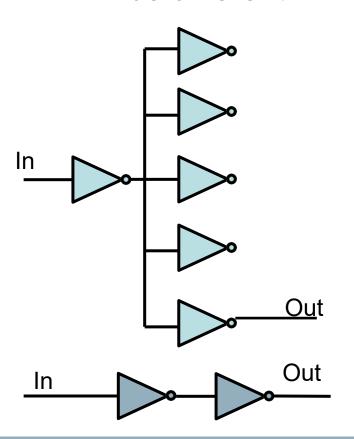
Capacitance can have a big **impact on transition times and delays.**

You will solve a similar exercise in tutorial 1

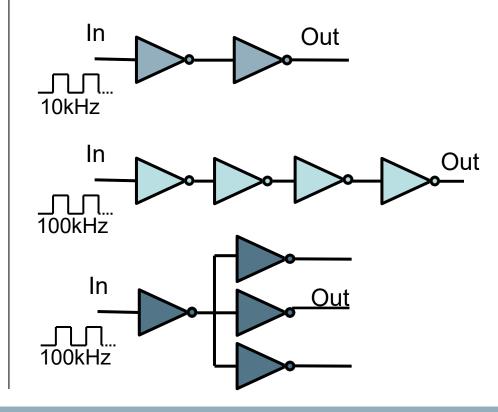


Visualising some concepts

Will the propagation times be different?



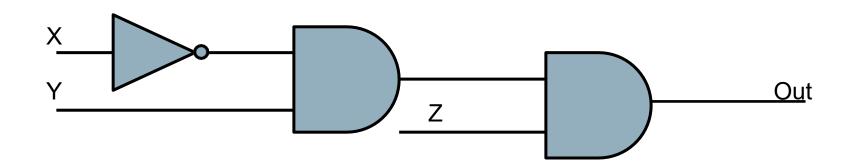
Which schematic is likely to produce the most heat?





Timing

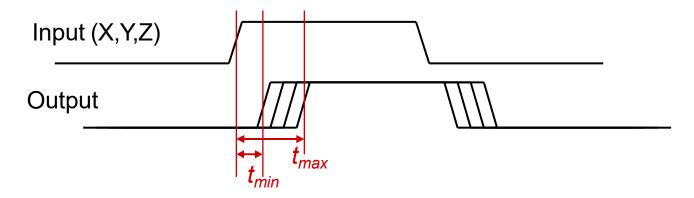
- We have seen that load and logic complexity can affect transition times
- It is common to have different propagation delays depending on the signal's pathway.
 - For example, in the circuit below, a change in input Z will show at the output after 1 gate has transitioned, while a change in Y will show after 2 transitions; X after 3 transitions.





Timing (2)

- This uncertainty in the switching timing of the output is reflected in timing diagrams. These are a graphical representation which can be useful to the designer in planning/refining a circuit design.
- Simulation softwares can be used to generate timing diagrams and you will learn how to do so during the labs.

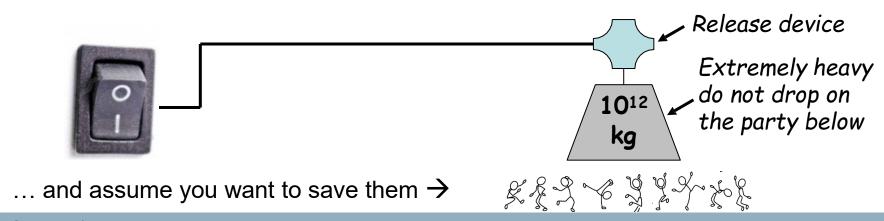


A good design should be expected to function for all input combinations.
 Therefore, considering worst-case or maximum delays is a safe way to ensure that the system will deliver dependable performance.



Hazards

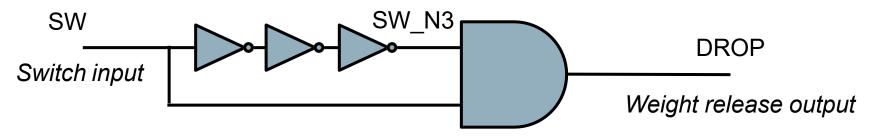
- In some cases paths with different propagation delays can create hazards. In purely combinational circuits, these should be avoided.
 - A hazard is a temporary "glitch" or incorrect output displayed as part
 of a transition in the circuit. Since combinational circuits (unlike sequential
 ones, which we will discuss shortly) react immediately to a change in input, a
 glitch could initiate a cascade of undesirable transitions. Causing the system
 to perform poorly.
 - For an example, look at this (silly) system ...



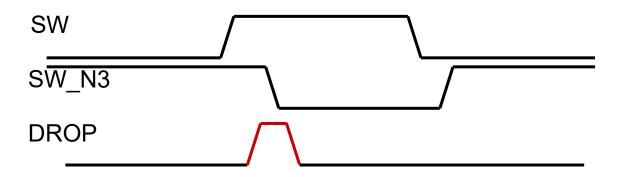


Hazards (2)

Would this system work?



You'd think the partygoers could breathe a sigh of relief, after all (SW·(((SW)')')') = 0. Right? Look at the timing diagram below.





Hazards (3)

- Although it might only last a few nanoseconds (transitions of 3 inverters), the glitch in the output *may* prove disastrous. **Steady state analysis is not indicative of dynamic performance.**
- This is a *static type 0 hazard*, where a change in the input which should see the output remain LOW determines a brief "1" glitch. There are also *type 1 hazards* (dual) and *dynamic hazards*.
 - See Wakerly 4.4
- Whether or not you need to worry about hazards depends on how critical the impact of the glitch is on the output.
- In sequential circuit design we can use **synchronisation** to neglect hazards. Signals are "sampled" at time intervals which are **far enough in time for all transient behaviour to have settled**.
 - You will see an example of this in CLAB 2



Summing up

- You now know about some non-ideal aspects of digital circuits with a particular focus on CMOS hardware:
 - Importance of correct voltages
 - For digital operation
 - To avoid excessive power dissipation
 - Role of loads in affecting timing
 - Hazards arising from propagation delays
- As you will learn to use computer-assisted design (CAD) tools, the software (ISE WebPACK in our lab activities) will assist you in diagnosing whether your design is at risk of undesirable behaviour.
 - Softwares are no substitute for a good understanding of the underlying principles. You need to have some idea of where to look when the CAD tools "complain" about your design.