Use the statement counting approach to determine the number of statements executed by the f---

### **Question 4**

Do the answers to question 3 change if the call to t.toString() is  $\Theta(n)$  rather than  $\Theta(1)$ ? Why or why not?

### Solution

Answer: No, the asymptotic complexities remain the same.

#### **Explanation:**

When analyzing asymptotic complexity, we consider how the dominant terms

#### Solution

#### 1. Initialization:

○ StringBuilder t = new StringBuilder(s); → 1 statement

#### 2. For Loop:

- Loop control (i < s.length()) and update (i++), plus the if condition check → 2 statements per iteration</li>
- Number of iterations: *n*
- One extra loop-condition check when the loop exits → +1 statement
- $\circ$  Total loop cost: 2n+1

#### 3. Return:

### Solution

Before we analyze, here is a quick reminder of what each notation means:

- **Big-O (O):** An upper bound on the growth rate of a function. It describes the *worst-case* scenario—how quickly the running time can grow as the input size increases.
- **Big-Omega** ( $\Omega$ ): A lower bound on the growth rate. It describes the *best-case* scenario—the minimum time the algorithm will take as the input size increases.
- **Big-Theta** (Θ): A tight bound. It means the function grows at the same rate in both the best and worst cases; the running time is always proportional to this rate.

#### Worst Case (Big-O):

In the worst case, every character in the string is a lowercase letter, so the body of the if statement executes every time. The loop runs for every character in the string, and all operations inside the loop are constant-time. Therefore, the worst-case time complexity is

#### • Best Case (Big-Ω):

In the best case, none of the characters are lowercase (so the if body never executes). However, the loop still iterates over every character, performing the if check each time. Thus, even in the best case, the time taken is proportional to the length of the string:

$$\Omega(n)$$

#### Overall Time Complexity (Big-Θ):

Since both the best case and worst case time complexities grow linearly with (

Express the overall time complexity of the method in question 5 using the appropriate notation(s) (big-O, big- $\Theta$ , big- $\Omega$ ).

### Solution

Answer: Θ(hw)

#### **Explanation:**

Since the best case and worst case are identical (as shown in Question 5), we can use Big- $\Theta$  notation to express a tight bound.

The exact statement count is 2hw + 2h + 5, where:

- The dominant term is **2hw** (quadratic in the dimensions)
- The linear term **2h** becomes negligible as h and w grow large

Express the overall time complexity of the method in question 7 using the appropriate notation(s) (big-O, big- $\Theta$ , big- $\Omega$ ).

### Solution

#### **Answer:**

- Worst case: O(n<sup>2</sup>)
- Best case:  $\Omega(1)$
- No tight bound (Θ) exists

#### **Explanation:**

The time complexity analysis reveals two fundamentally different execution scenarios:

Determine the time complexity of the following pseudocode snippet using the active operation approach (careful! This one is tricky!).

```
Let G be a weighted graph implemented with an adjacency list;
Let H be a heap of edges ordered by edge weight (initially empty);

for (Vertex v : G.vertices()) {
    if (G.hasEdge(v, 0)) {
        H.insert(new Edge(v, 0));
    }
}
```

#### Solution

### Why This is Tricky

This problem requires careful analysis because two operations have non-

Determine exact number of statements executed by the following pseudocode function in the worst case (this one is a bit tricky too!):

```
public String treeConcat(Node root) {
   String s = root.string();

   if (!root.isLeaf()) {
      for (Node child : root.children()) {
            s += treeConcat(child);
        }
   }

   return s;
}
```

### Solution