

# ICS Homework 12

May 31, 2020

## 1 Organization

### 1.1

If we only use 0x400000-0x800000, compare the memory usage of page table between single level and 4 level.

Single:  $\frac{2^{48}B}{4KB} * 8B = 2^{39}B = 512GB$

4 level:  $(2 \text{ L4 page table} + 3 \text{ L1/L2/L3 page table}) * 4KB = 20KB$

### 1.2

If we have enough physical memory and use the whole  $2^{48}B$  virtual space, compare the memory usage again.

Single:  $(2^{48}B/4KB) * 8B = 2^{39}B = 512GB$

4 level:  $(1 + 512 + 512^2 + 512^3) * 4KB \approx 513GB$

### 1.3

Given the following page table and cr3=0x1000, please translate the virtual addresses to physical addresses. Assume the machine is x86\_64, which uses 4 level page table.

| Physical Address of PTE | Address Part in PTE |
|-------------------------|---------------------|
| 0x1000                  | 0x4000              |
| 0x1008                  | 0x4000              |
| 0x4000                  | 0x5000              |
| 0x5000                  | 0x7000              |
| 0x5008                  | 0x8000              |
| 0x5010                  | 0x9000              |
| 0x5018                  | 0x3000              |
| 0x5020                  | 0x5000              |
| 0x7008                  | 0x8000              |
| 0x7010                  | 0x10000             |

| Physical Address | Virtual Address |
|------------------|-----------------|
| 0x1234           | 0x8234          |
| 0x2234           | 0x10234         |
| 0x8000802135     | 0x9135          |

## 1.4

Assume on a x86\_64 machine without cache, we have an int `a[70]` at virtual address `0x8000344f00`, a program access `a[0]`, `a[1]`, ..., `a[69]` one by one. Before the first access, we have only an empty L1 page table.

### 1.4.1

How many physical page is used after the program access `a[0]`, `a[64]`, `a[69]`? (Remember the page table is stored in physical page too)

after access `a[0]`: 4 page table pages + 1 content pages

after access `a[64]`: 4 + 2

after access `a[69]`: 4 + 2

### 1.4.2

How many memory accesses and page fault happened in the program if we have no TLB?

$(4+1)*70+1+4=355$ , 4+1 for each successful array access.

First access to `a[0]` need one access to L1 PTE to trigger page fault. First access to `a[64]` need four access to L1, L2, L3 and L4 PTE to trigger page fault. 2 page faults in total.

### 1.4.3

How many memory accesses and page fault happened in the program if we have a full associative infinite TLB?

$70+4*2+1+4=83$ . 1 for each successful array access. Accesses to `a[0]`, `a[64]` will trigger two page table walk and fill TLB.

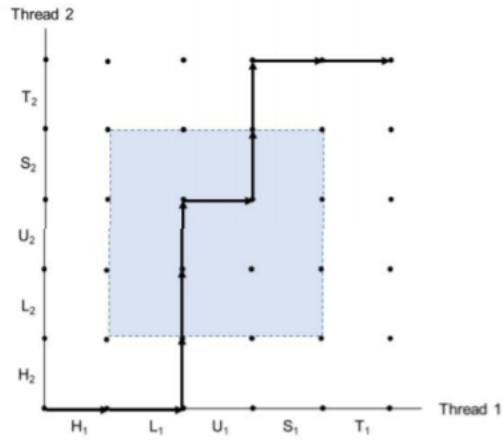
2 page faults in total.

## 2 System Software

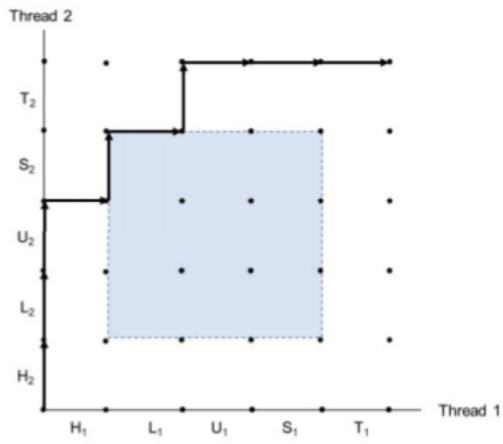
### 2.1

Using the progress graph in Figure 12-21 of file "badcnt.c", draw the following trajectories out and point out the value of `cnt` after the execution (assume the value of `cnt` is 0 initially for each trajectory).

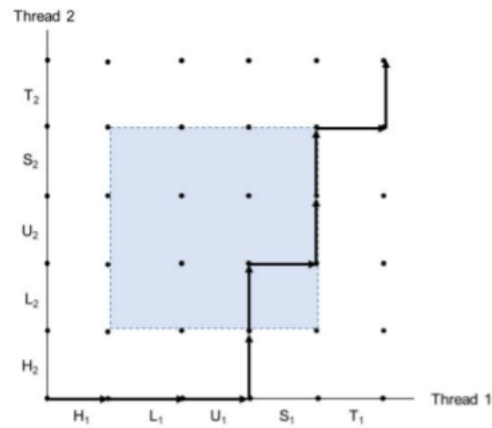
1. H1,L1,H2,L2,U2,U1,S2,T2,S1,T1  
cnt = 1



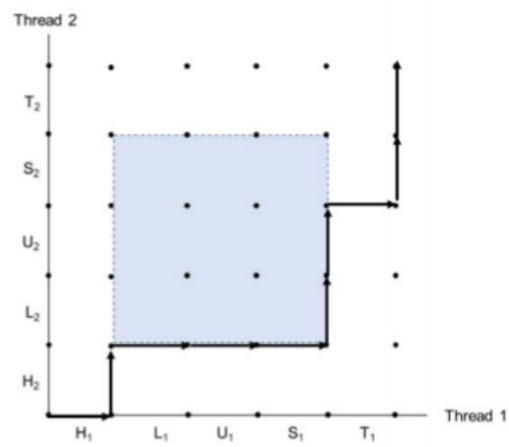
2.  $H_2, L_2, U_2, H_1, S_2, L_1, T_2, U_1, S_1, T_1$   
 $\text{cnt} = 2$



3.  $H_1, L_1, U_1, H_2, L_2, S_1, U_2, S_2, T_1, T_2$   
 $\text{cnt} = 1$



4. H1,H2,L1,U1,S1,L2,U2,T1,S2,T2  
cnt = 2



## 2.2

1. Complete the previous code according to the comment.

```

1  #include "csapp.h"
2  #define N 4
3
4  void *thread(void *vargp) {
5      int myid = *((int)vargp);
6      printf("in thread %d\n", myid);
7      return NULL;
8  }
9
10 int main() {

```

```

11 pthread_t tid[N];
12 int *ptr;
13
14 for (int i = 0; i < N; i++) {
15     ptr = malloc(sizeof(int));
16     *ptr = i;
17     // create a thread running
18     // with argument ptr
19     // your code here
20     pthread_create(&tid[i], NULL,
21                  thread, ptr);
22
23     free(ptr);
24 }
25
26 for (int i = 0; i < N; i++)
27     pthread_join(tid[i], NULL);
28 }

```

2. Is there any race condition in the previous code? Why or why not?

SOL:

Yes. If the free call executed before the newly created thread, then there will be a segmentation fault caused by accessing a freed pointer.