Algorithms and Data Structures Graphs, Depth-/Breadth-first Search, Graph-Connectivity

Albert-Ludwigs-Universität Freiburg

Prof. Dr. Rolf Backofen

Bioinformatics Group / Department of Computer Science Algorithms and Data Structures, January 2019

Structure



Graphs

Introduction Implementation Application example



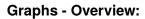


Besides arrays, lists and trees the most common data structure (Trees are a special type of graph)

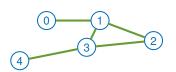
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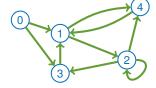
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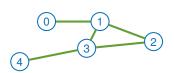
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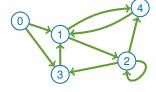


- Besides arrays, lists and trees the most common data structure
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- Breadth-first search (BFS)
- Depth-first search (DFS)
- Connected components of a graph

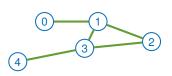


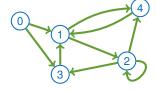




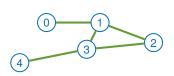


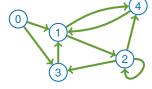
■ Each graph G = (V, E) consists of:





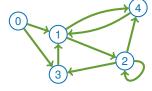
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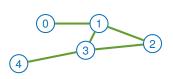


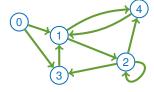
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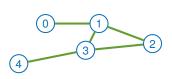


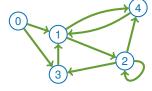
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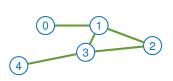


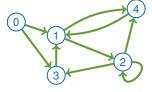
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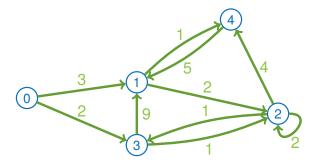




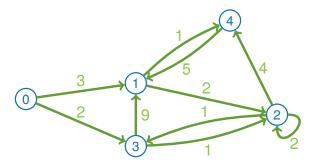
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 - Undirected edge: $e = \{u, v\}$ (set)
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- Self-loops are also possible: e = (u, u) or $e = \{u, u\}$

Weighted graph:

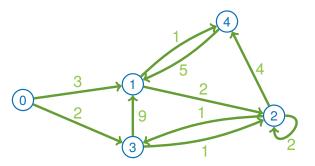
Weighted graph:



Weighted graph:



Each edge is marked with a real number named weight



- Each edge is marked with a real number named weight
- The weight is also named length or cost of the edge depending on the application

Graphs Introduction



Example: Road network



Example: Road network

Intersections:

vertices



Example: Road network

Intersections:

vertices

■ Roads: edges



Example: Road network

Intersections:

vertices

■ Roads: edges

Travel time:

costs of the edges



Intersections: vertices

■ Roads: edges

Travel time: costs of the edges



Figure: Map of Freiburg © OpenStreetMap

Structure



Graphs

Introduction

Implementation

Application example

Adjacency matrix with space consumption $\Theta(|V|^2)$

11 Adjacency matrix with space consumption $\Theta(|V|^2)$

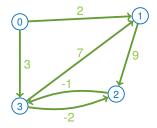


Figure: Weighted graph with

$$|V| = 4$$
, $|E| = 6$

Adjacency matrix with space consumption $\Theta(|V|^2)$

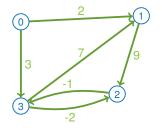


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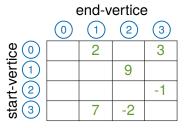


Figure: Adjacency matrix



2 Adjacency list / fields with space consumption $\Theta(|V| + |E|)$

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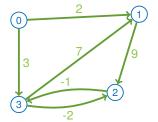


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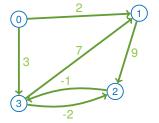


Figure: Weighted graph with |V| = 4, |E| = 6

<u>o</u>	1, 2	3, 3
ert (1)	2, 9	
start-vertice	3, -1	
sta 3	1, 7	2, -2

Figure: Adjacency list

Graphs Implementation

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Graphs

Implementation



Graph: Arrangement

■ Graph is fully defined through the adjacency matrix / list

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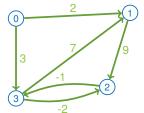


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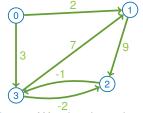


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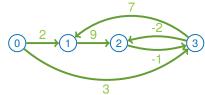


Figure: Same graph ordered by number - outer planar graph



```
class Graph:
    def __init__(self):
        self.vertices = []
        self.edges = []
    def addVertice(self, vert):
        self.vertices.append(vert)
    def addEdge(self, fromVert, toVert, cost):
        self.edges.append( \
            (fromVert, toVert, cost))
```

Graphs

Degrees (Valency)

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Degree of a vertex: Directed graph: G = (V, E)

Degrees (Valency)

Degree of a vertex: Directed graph: G = (V, E)



Figure: Vertex with in- / outdegree of 3 / 2

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Indegree of a vertex u is the number of edge head ends adjacent to the vertex

$$\deg^+(u) = |\{(v, u) : (v, u) \in E\}|$$



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Outdegree of a vertex u is the number of edge tail ends adjacent to the vertex

$$\deg^-(\mathbf{u}) = |\{(\mathbf{u}, \mathbf{v}) : (\mathbf{u}, \mathbf{v}) \in \mathbf{E}\}|$$

Graphs

Degrees (Valency)



Degree of a vertex: Undirected graph: G = (V, E)

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Figure: Vertex with degree of 4

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Degree of a vertex u is the number of vertices adjacent to the vertex

$$\deg(u) = |\{\{v, u\} : \{v, u\} \in E\}|$$

Graphs Paths



FREE

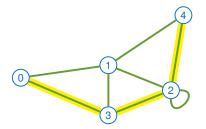


Figure: Undirected path of length 3 P = (0,3,2,4)

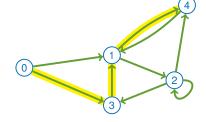


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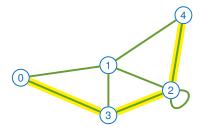


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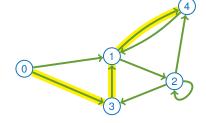


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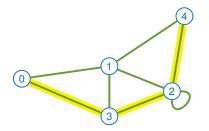


Figure: Undirected path of length 3 P = (0, 3, 2, 4)

Figure: Directed path of length 3 P = (0.3, 1.4)

- A path of G is a sequence of edges $u_1, u_2, ..., u_i \in V$ with
 - Undirected graph: $\{u_1, u_2\}, \{u_2, u_3\}, \dots, \{u_{i-1}, u_i\} \in E$
 - Directed graph: $(u_1, u_2), (u_2, u_3), \dots, (u_{i-1}, u_i) \in E$

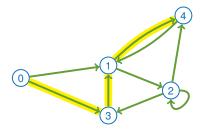


Figure: Directed path of length 3 P = (0, 3, 1, 4)

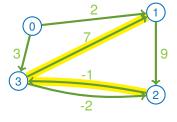


Figure: Weighted path with cost 6 P = (2, 3, 1)

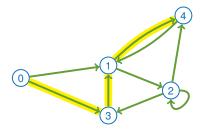


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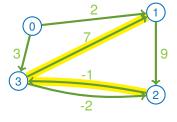


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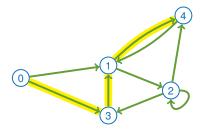


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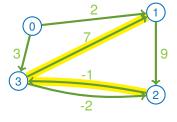
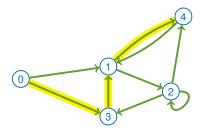


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Paths in a graph: G = (V, E)



3 -1 2

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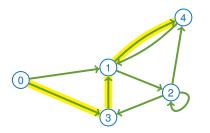


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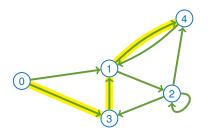


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 - With weights: sum of weigths of edges taken

Shortest path in a graph: G = (V, E)

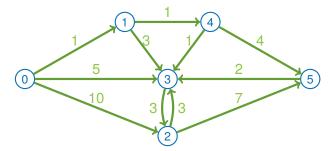


Figure: Shortest path from 0 to 2 with cost / distance d(0,2) = ?

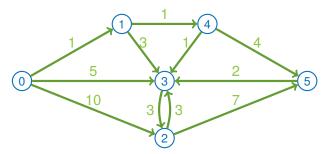


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The shortest path between two vertices u, v is the path P = (u, ..., v) with the shortest length d(u, v) or lowest costs

Shortest path in a graph: G = (V, E)

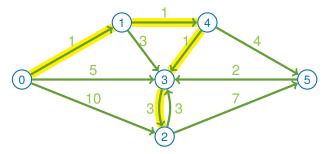


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Graphs Paths



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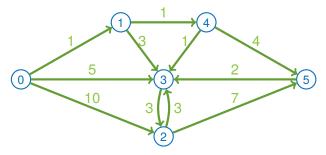


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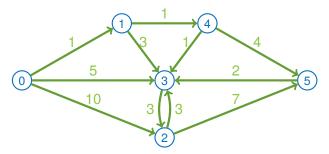


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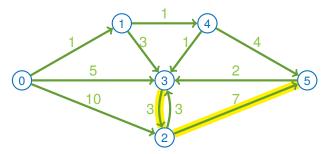


Figure: Diameter of graph is d = 10, P = (3, 2, 5)

The diameter of a graph is the length / the costs of the longest shortest path

Graphs Connected Components

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Connected components: G = (V, E)

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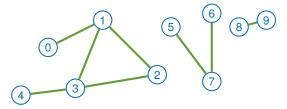


Figure: Three connected components

Undirected graph:

Connected components: G = (V, E)

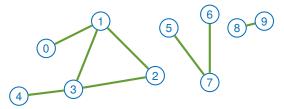


Figure: Three connected components

- Undirected graph:
 - All connected components are a partition of V

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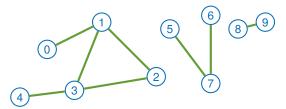


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- Undirected graph:
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Graphs Connected Components



Connected components: G = (V, E)

Graphs Connected Components



Connected components: G = (V, E)

Directed graph:

Connected components: G = (V, E)

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Connected components: G = (V, E)

- Directed graph:
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 - Not part of this lecture

Graphs

Connected Components - Graph Exploration



Graph Exploration: (Informal definition)

Graphs

Connected Components - Graph Exploration



Graph Exploration: (Informal definition)

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- We visit each reachable vertex connected to s

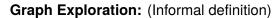
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 - Searching of connected components
 - Flood fill in drawing programms

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- Mark the start vertex s (level 0)

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- Mark all unmarked connected vertices (level 1)
- Mark all unmarked vertices connected to a level 1-vertex (level 2)
- 5 Iteratively mark reachable vertices for all levels

- We start with all vertices unmarked and mark visited vertices
- Mark the start vertex s (level 0)
- Mark all unmarked connected vertices (level 1)
- Mark all unmarked vertices connected to a level 1-vertex (level 2)
- 5 Iteratively mark reachable vertices for all levels
- All connected nodes are now marked and in the same connected component as the start vertex s

Graphs

Connected Components - Breadth-First Search



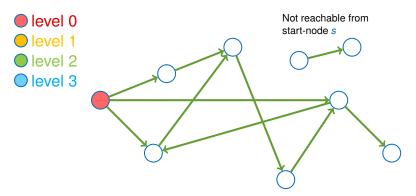


Figure: spanning tree of a breadth-first search

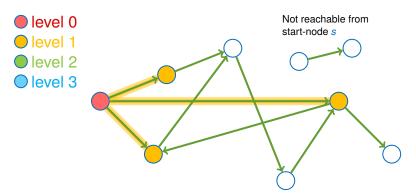


Figure: spanning tree of a breadth-first search

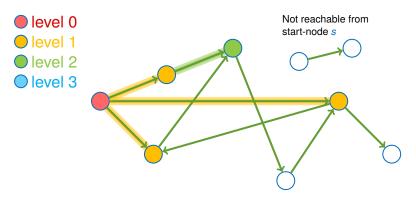


Figure: spanning tree of a breadth-first search

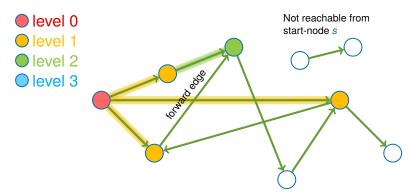


Figure: spanning tree of a breadth-first search

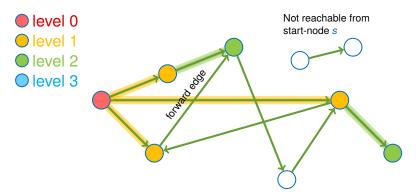


Figure: spanning tree of a breadth-first search

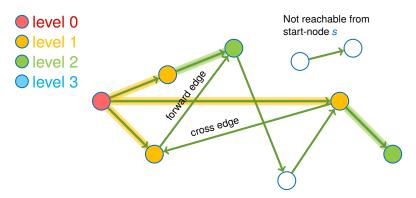


Figure: spanning tree of a breadth-first search

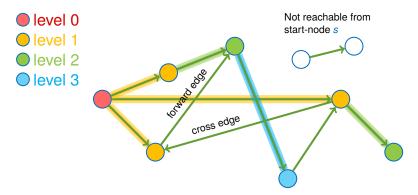


Figure: spanning tree of a breadth-first search

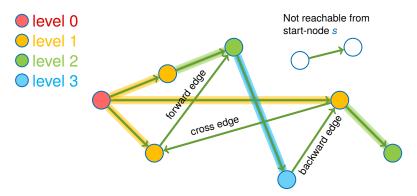


Figure: spanning tree of a breadth-first search

We start with all vertices unmarked and mark visited vertices

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- Pick an unmarked connected vertex and start a recursive depth-first search with the vertex as start vertex (continue on step 2)
- If no unmarked connected vertex exists go one vertex back and continue recursive search (reduce the recursion level by one)

Search starts with long paths (searching with depth)

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 - Each newly visited vertex gets marked by an increasing number

Depth-first search:

- Search starts with long paths (searching with depth)
- Like breadth-first search marks all connected vertices
- If the graph is acyclic we get a topological sorting
 - Each newly visited vertex gets marked by an increasing number
 - The numbers increase with path length from the start vertex

Graphs

Connected Components - Depth-First Search

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- The marked vertices create a different spanning tree containing all reachable nodes
- start-node
- path 1
- path 2
- opath 3

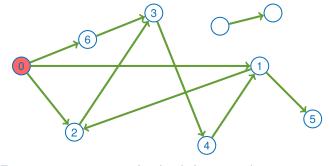


Figure: spanning tree of a depth-first search

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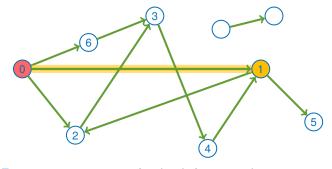


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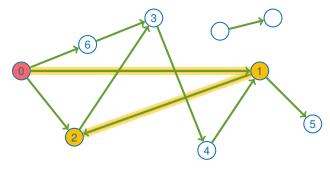


Figure: spanning tree of a depth-first search

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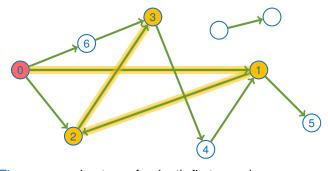


Figure: spanning tree of a depth-first search

- start-node
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- path 2
- opath 3

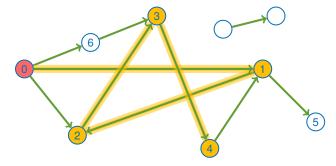


Figure: spanning tree of a depth-first search

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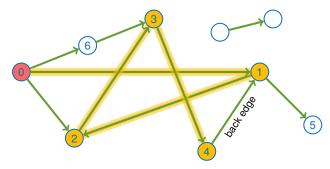


Figure: spanning tree of a depth-first search

- start-node
- path 1
- path 2
- path 3

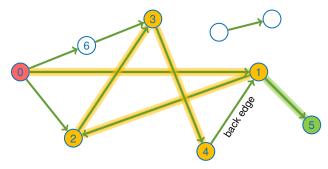


Figure: spanning tree of a depth-first search

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- opath 3

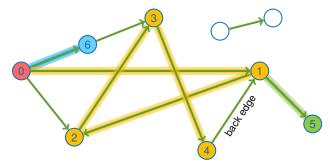


Figure: spanning tree of a depth-first search

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- opath 3

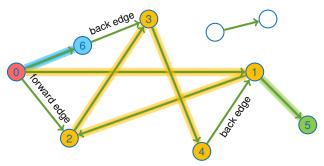


Figure: spanning tree of a depth-first search

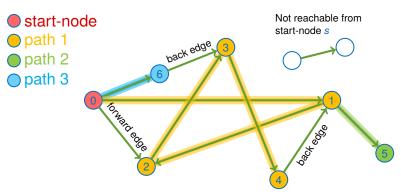


Figure: spanning tree of a depth-first search

Graphs

Why is this called Breadth- and Depth-First Search?



Constant costs for each visited vertex and edge

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- We get a runtime complexity of $\Theta(|V'| + |E'|)$
- Let V' and E' be the reachable vertices and edges
- All vertices of V' are in the same connected component as our start vertex s
- This can only be improved by a constant factor



Graphs

Introduction Implementation

Application example

Image processing



Image processing



■ Connected component labeling

Image processing

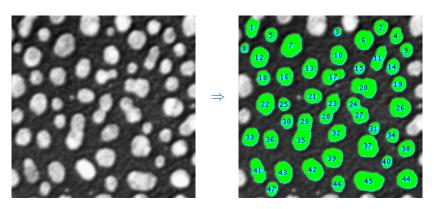


- Connected component labeling
- Counting of objects in an image

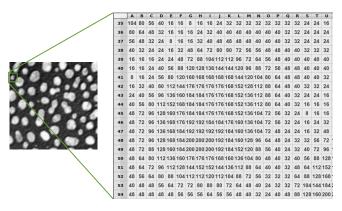
Image processing



- Connected component labeling
- Counting of objects in an image

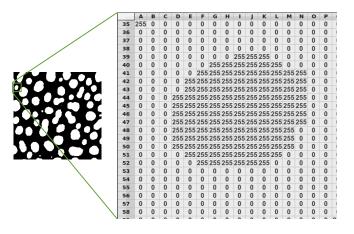


What is object, what is background?



Convert to black and white using threshold:

value = 255 if value > 100 else 0



Application example Image processing



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Interpret image as graph:

Application example Image processing

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Each white pixel is a node

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- Edges between adjacent pixels (normally 4 or 8 neighbors)

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Interpret image as graph:

- Each white pixel is a node
- Edges between adjacent pixels (normally 4 or 8 neighbors)
- Edges are not saved externally, algorithm works directly on array
- Breadth- / depth-first search find all connected components (particles)

Image processing



Find connected components:

Image processing



Find connected components:

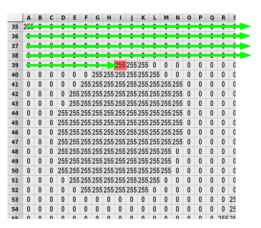
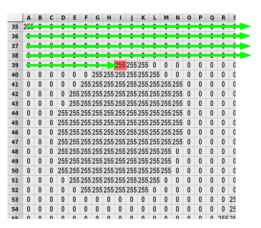


Image processing

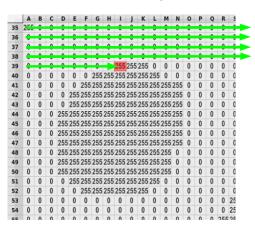


Find connected components:



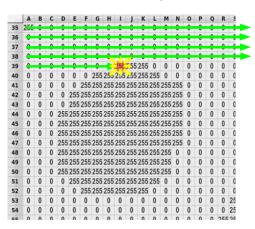
Search pixel-by-pixel for non-zero intensity

Find connected components:



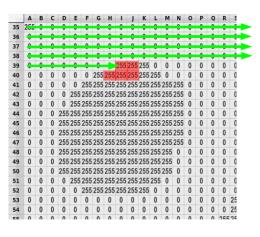
- Search pixel-by-pixel for non-zero intensity
- Label found pixel as component 1

Find connected components:



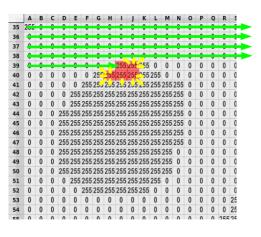
- Search pixel-by-pixel for non-zero intensity
- Label found pixel as component 1
- Check neighbors of all new labeled pixels



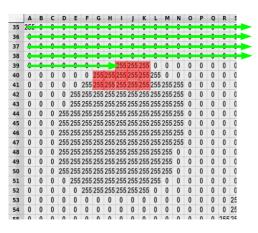


- Search pixel-by-pixel for non-zero intensity
- Label found pixel as component 1
- Check neighbors of all new labeled pixels
- Label non-zero pixels as component 1



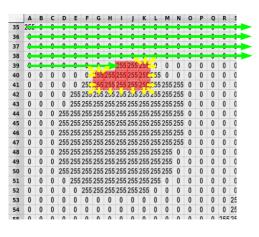


- Search pixel-by-pixel for non-zero intensity
- Label found pixel as component 1
- Check neighbors of all new labeled pixels
- Label non-zero pixels as component 1



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- Label found pixel as component 1
- Check neighbors of all new labeled pixels
- Label non-zero pixels as component 1



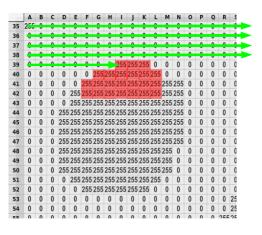


- Search pixel-by-pixel for non-zero intensity
- Label found pixel as component 1
- Check neighbors of all new labeled pixels
- Label non-zero pixels as component 1

Application example

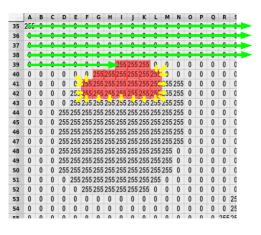
Image processing





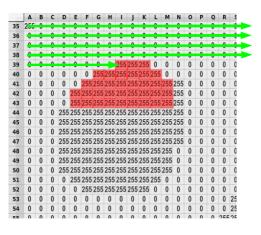
- Search pixel-by-pixel for non-zero intensity
- Label found pixel as component 1
- Check neighbors of all new labeled pixels
- Label non-zero pixels as component 1





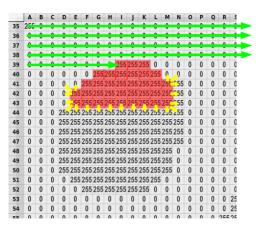
- Search pixel-by-pixel for non-zero intensity
- Label found pixel as component 1
- Check neighbors of all new labeled pixels
- Label non-zero pixels as component 1





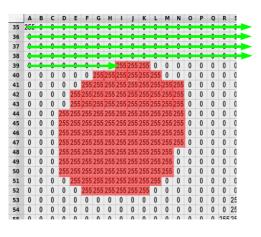
- Search pixel-by-pixel for non-zero intensity
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- Label non-zero pixels as component 1





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- Label found pixel as component 1
- Check neighbors of all new labeled pixels
- Label non-zero pixels as component 1



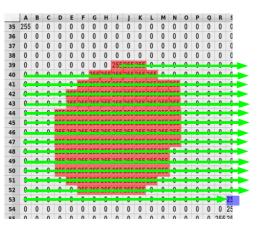


- Search pixel-by-pixel for non-zero intensity
- Label found pixel as component 1
- Check neighbors of all new labeled pixels
- Label non-zero pixels as component 1

Application example

Image processing





- Search pixel-by-pixel for non-zero intensity
- Label found pixel as component 2
- ...

Result of connected component labeling:

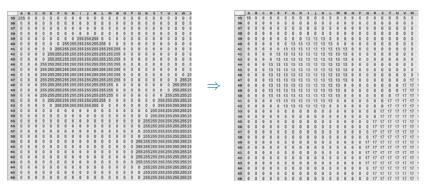


Figure: Result: particle indices instead of intensities

■ General

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■ Graph Search

Graph Connectivity

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[Wik] Connectivity (graph theory)
    https://en.wikipedia.org/wiki/Connectivity_
        (graph_theory)
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